CP/M 2

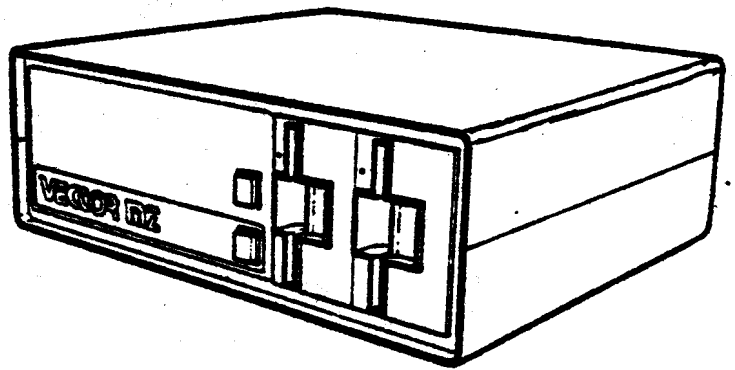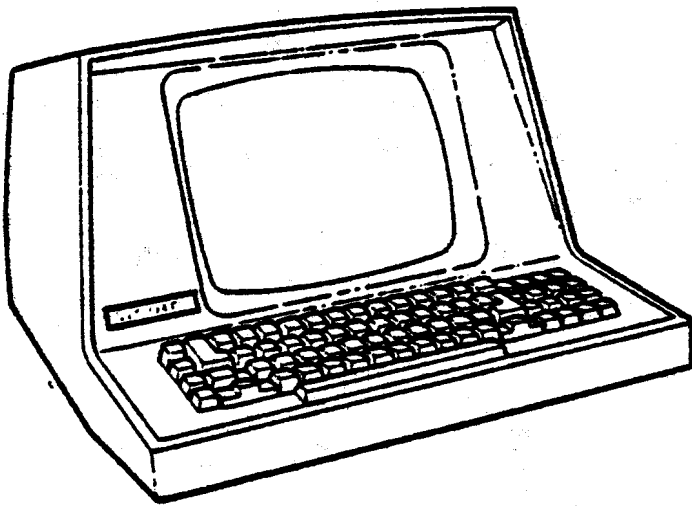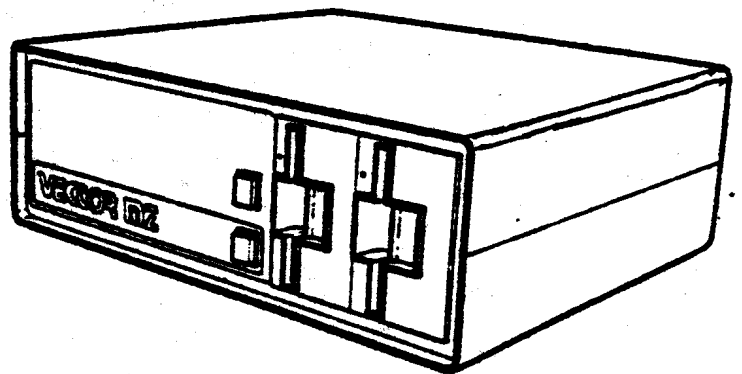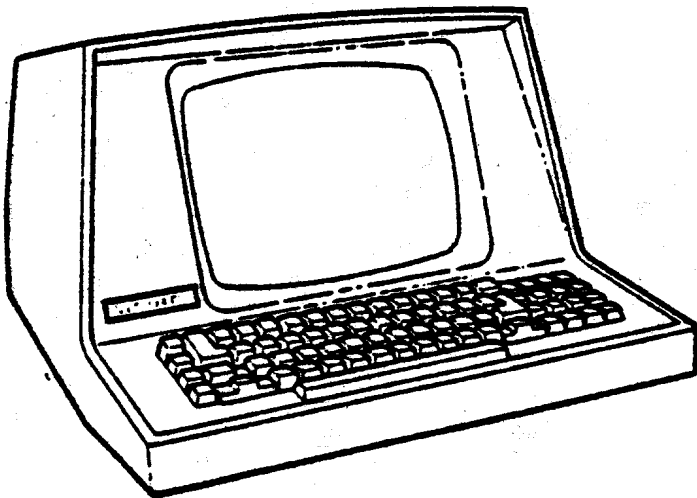## USERS GUIDE FOR 1.4 OWNERS

VECTOR
VECTOR GRAPHIC, INC.

CP/M 2

USERS GUIDE FOR 1.4 OWNERS

VECTOR
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CP/M USER'S GUIDE

FOR CP/M 1.4 OWNERS

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CP/M 2.0 USER'S GUIDE FOR CP/M 1.4 OWNERS

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1. AN OVERVIEW OF CP/M 2.0 FACILITIES.

CP/M 2.0 is a high-performance single-console operating system which uses table driven techniques to allow field reconfiguration to match a wide variety of disk capacities. All of the fundamental file restrictions are removed, while maintaining upward compatibility from previous versions of release 1. Features of CP/M 2.0 include field specification of one to sixteen logical drives, each containing up to eight megabytes. Any particular file can reach the full drive size with the capability to expand to thirty-two megabytes in future releases. The directory size can be field configured to contain any reasonable number of entries, and each file is optionally tagged with read/only and system attributes. Users of CP/M 2.0 are physically separated by user numbers, with facilities for file copy operations from one user area to another. Powerful relative-record random access functions are present in CP/M 2.0 which provide direct access to any of the 65536 records of an eight megabyte file.

All disk-dependent portions of CP/M 2.0 are placed into a BIOS-resident "disk parameter block" which is either hand coded or produced automatically using the disk definition macro library provided with CP/M 2.0. The end user need only specify the maximum number of active disks, the starting and ending sector numbers, the data allocation size, the maximum extent of the logical disk, directory size information, and reserved track values. The macros use this information to generate the appropriate tables and table references for use during CP/M 2.0 operation. Deblocking information is also provided which aids in assembly or disassembly of sector sizes which are multiples of the fundamental 128 byte data unit, and the system alteration manual includes general-purpose subroutines which use the this deblocking information to take advantage of larger sector sizes. Use of these subroutines, together with the table driven data access algorithms, make CP/M 2.0 truly a universal data management system.

File expansion is achieved by providing up to 512 logical file extents, where each logical extent contains 16K bytes of data. CP/M 2.0 is structured, however, so that as much as 128K bytes of data is addressed by a single physical extent (corresponding to a single directory entry), thus maintaining compatibility with previous versions while taking full advantage of directory space.

Random access facilities are present in CP/M 2.0 which allow immediate reference to any record of an eight megabyte file. Using CP/M's unique data organization, data blocks are only allocated when actually required and movement to a record position requires little search time. Sequential file access is upward compatible from earlier versions to the full eight megabytes, while random access compatibility stops at 512K byte files. Due to CP/M 2.0's simpler and faster random access, application programmers are encouraged to alter their programs to take full advantage of the 2.0 facilities.

Several CP/M 2.0 modules and utilities have improvements which correspond to the enhanced file system. STAT and PIP both account for file attributes and user areas, while the CCP provides a "login"

function to change from one user area to another. The CCP also formats directory displays in a more convenient manner and accounts for both CRT and hard-copy devices in its enhanced line editing functions.

The sections below point out the individual differences between CP/M 1.4 and CP/M 2.0, with the understanding that the reader is either familiar with CP/M 1.4, or has access to the 1.4 manuals.

2. USER INTERFACE.

Console line processing takes CRT-type devices into account with three new control characters, shown with an asterisk in the list below (the symbol "ctl" below indicates that the control key is simultaneously depressed):

```
rub/del removes and echoes last character
  ctl-C  reboot when at beginning of line
  ctl-E  physical end of line
  ctl-H  backspace one character position*
  ctl-J  (line feed) terminates current input*
  ctl-M  (carriage return) terminates input
  ctl-R  retype current line after new line
  ctl-U  remove current line after new line
  ctl-X  backspace to beginning of current line*
```

In particular, note that ctl-H produces the proper backspace overwrite function (ctl-H can be changed internally to another character, such as delete, through a simple single byte change). Further, the line editor keeps track of the current prompt column position so that the operator can properly align data input following a ctl-U, ctl-R, or ctl-X command.

3. CONSOLE COMMAND PROCESSOR (CCP) INTERFACE.

There are four functional differences between CP/M 1.4 and CP/M 2.0 at the console command processor (CCP) level. The CCP now displays directory information across the screen (four elements per line), the USER command is present to allow maintenance of separate files in the same directory, and the actions of the "ERA *.*" and "SAVE" commands have changed. The altered DIR format is self-explanatory, while the USER command takes the form:

USER n

where n is an integer value in the range 0 to 15. Upon cold start, the operator is automatically "logged" into user area number 0, which is compatible with standard CP/M 1.4 directories. The operator may issue the USER command at any time to move to another logical area within the same directory. Drives which are logged-in while addressing one user number are automatically active when the operator moves to another user number since a user number is simply a prefix which accesses particular directory entries on the active disks.

The active user number is maintained until changed by a subsequent USER command, or until a cold start operation when user 0 is again assumed.

Due to the fact that user numbers now tag individual directory entries, the ERA *.* command has a different effect. In version 1.4, this command can be used to erase a directory which has "garbage" information, perhaps resulting from use of a diskette under another operating system (heaven forbid!). In 2.0, however, the ERA *.* command affects only the current user number. Thus, it is necessary to write a simple utility to erase a nonsense disk (the program simply writes the hexadecimal pattern E5 throughout the disk).

The SAVE command in version 1.4 allows only a single memory save operation, with the potential of destroying the memory image due to directory operations following extent boundary changes. Version 2.0, however, does not perform directory operations in user data areas after disk writes, and thus the SAVE operation can be used any number of times without altering the memory image.

4. STAT ENHANCEMENTS.

The STAT program has a number of additional functions which allow disk parameter display, user number display, and file indicator manipulation. The command:

STAT VAL:

produces a summary of the available status commands, resulting in the output:

```
        Temp R/O Disk: d:=R/O
        Set Indicator: d:filename.typ $R/O $R/W $SYS $DIR
        Disk Status  : DSK: d:DSK:
        User Status  : USR:
        Iobyte Assign:
        (list of possible assignments)
```

which gives an instant summary of the possible STAT commands. The command form:

STAT d:filename.typ $S

where "d:" is an optional drive name, and "filename.typ" is an unambiguous or ambiguous file name, produces the output display format:

```
        Size   Recs  Bytes  Ext Acc
          48     48     6k    1 R/O A:ED.COM
          55     55    12k    1 R/O (A:PIP.COM)
       65536    128     2k    2 R/W A:X.DAT
```

where the $S parameter causes the "Size" field to be displayed (without the $S, the Size field is skipped, but the remaining fields are displayed). The Size field lists the virtual file size in records, while the "Recs" field sums the number of virtual records in each extent. For files constructed sequentially, the Size and Recs fields are identical. The "Bytes" field lists the actual number of bytes allocated to the corresponding file. The minimum allocation unit is determined at configuration time, and thus the number of bytes corresponds to the record count plus the remaining unused space in the last allocated block for sequential files. Random access files are given data areas only when written, so the Bytes field contains the only accurate allocation figure. In the case of random access, the Size field gives the logical end-of-file record position and the Recs field counts the logical records of each extent (each of these extents, however, may contain unallocated "holes" even though they are added into the record count). The "Ext" field counts the number of logical 16K extents allocated to the file. Unlike version 1.4, the Ext count does not necessarily correspond to the number of directory entries given to the file, since there can be up to 128K bytes (8 logical extents) directly addressed by a single directory entry, depending upon allocation size (in a special case, there are actually 256K bytes which can be directly addressed by a physical extent).

The "Acc" field gives the R/O or R/W access mode, which is changed using the commands shown below. Similarly, the parentheses

shown around the PIP.COM file name indicate that it has the "system" indicator set, so that it will not be listed in DIR commands. The four command forms

```
STAT d:filename.typ $R/O
STAT d:filename.typ $R/W
STAT d:filename.typ $SYS
STAT d:filename.typ $DIR
```

set or reset various permanent file indicators.   The R/O indicator places the file (or set of files) in a read-only status until changed by a subsequent STAT command. The R/O status is recorded in the directory with the file so that it remains R/O through intervening cold start operations.   The R/W indicator places the file in a permanent read/write status. The SYS indicator attaches the system indicator to the file, while the DIR command removes the system indicator. The "filename.typ" may be ambiguous or unambiguous, but in either case, the files whose attributes are changed are listed at the console when the change occurs. The drive name denoted by "d:" is optional.

When a file is marked R/O, subsequent attempts to erase or write into the file result in a terminal BDOS message

                    Bdos Err on d: File R/O

The BDOS then waits for a console input before performing a subsequent warm start (a "return" is sufficient to continue). The command form

                         STAT d:DSK:

lists the drive characteristics of the disk named by "d:" which is in the range A:, B:, ..., P:. The drive characteristics are listed in the format:

```
      d: Drive Characteristics
  65536: 128 Byte record Capacity
   8192: Kilobyte Drive Capacity
    128: 32 Byte Directory Entries
      0: Checked  Directory Entries
   1024: Records/ Extent
    128: Records/ Block
     58: Sectors/ Track
      2: Reserved Tracks
```

where "d:" is the selected drive, followed by the total record capacity (65536 is an 8 megabyte drive), followed by the total capacity listed in Kilobytes. The directory size is listed next, followed by the "checked" entries. The number of checked entries is usually identical to the directory size for removable media, since this mechanism is used to detect changed media during CP/M operation without an intervening warm start. For fixed media, the number is usually zero, since the media is not changed without at least a cold or warm start. The number of records per extent determines the addressing capacity of each directory entry (1024 times 128 bytes, or

128K in the example above). The number of records per block shows the basic allocation size (in the example, 128 records/block times 128 bytes per record, or 16K bytes per block). The listing is then followed by the number of physical sectors per track and the number of reserved tracks. For logical drives which share the same physical disk, the number of reserved tracks may be quite large, since this mechanism is used to skip lower-numbered disk areas allocated to other logical disks. The command form

STAT DSK:

produces a drive characteristics table for all currently active drives. The final STAT command form is

STAT USR:

which produces a list of the user numbers which have files on the currently addressed disk. The display format is:

        Active User : 0
        Active Files: 0 1 3

where the first line lists the currently addressed user number, as set by the last CCP USER command, followed by a list of user numbers scanned from the current directory. In the above case, the active user number is 0 (default at cold start), with three user numbers which have active files on the current disk. The operator can subsequently examine the directories of the other user numbers by logging-in with USER 1, USER 2, or USER 3 commands, followed by a DIR command at the CCP level.

5. PIP ENHANCEMENTS.

PIP provides three new functions which account for the features of CP/M 2.0. All three functions take the form of file parameters which are enclosed in square brackets following the appropriate file names. The commands are:

Gn	Get File from User number n
	(n in the range 0 - 15)
W·	Write over R/O files without
	console interrogation
R	Read system files

The G command allows one user area to receive data files from another. Assuming the operator has issued the USER 4 command at the CCP level, the PIP statement

                    PIP X.Y = X.Y[G2]

reads file X.Y from user number 2 into user area number 4. The command
                    PIP A:=A:*.*[G2]

copies all of the files from the A drive directory for user number 2 into the A drive directory of the currently logged user number. Note that to ensure file security, one cannot copy files into a different area than the one which is currently addressed by the USER command.

Note also that the PIP program itself is initially copied to a user area (so that subsequent files can be copied) using the SAVE command. The sequence of operations shown below effectively moves PIP from one user area to the next.

USER 0	login user 0
DDT PIP.COM	load PIP to memory
(note PIP size s)	
G0	return to CCP
USER 3	login user 3
SAVE s PIP.COM	

where s is the integral number of memory "pages" (256 byte segments) occupied by PIP. The number s can be determined when PIP.COM is loaded under DDT, by referring to the value under the "NEXT" display. If for example, the next available address is 1D00, then PIP.COM requires 1C hexadecimal pages (or 1 times 16 + 12 = 28 pages), and thus the value of s is 28 in the subsequent save. Once PIP is copied in this manner, it can then be copied to another disk belonging to the same user number through normal pip transfers.

Under normal operation, PIP will not overwrite a file which is set to a permanent R/O status. If attempt is made to overwrite a R/O file, the prompt

DESTINATION FILE IS R/O, DELETE (Y/N)?

is issued. If the operator responds with the character "y" then the
file is overwritten. Otherwise, the response

** NOT DELETED **

is issued, the file transfer is skippped, and PIP continues with the
next operation in sequence. In order to avoid the prompt and response
in the case of R/O file overwrite, the command line can include the W
parameter, as shown below

PIP A:=B:*.COM[W]

which copies all non-system files to the A drive from the B drive, and
overwrites any R/O files in the process. If the operation involves
several concatenated files, the W parameter need only be included with
the last file in the list, as shown in the following example

PIP A.DAT = B.DAT,F:NEW.DAT,G:OLD.DAT[W]

     Files with the system attribute can be included in PIP transfers
if the R parameter is included, otherwise system files are not
recognized. The command line

PIP ED.COM = B:ED.COM[R]

for example, reads the ED.COM file from the B drive, even if it has
been marked as a R/O and system file. The system file attributes are
copied, if present.

     It should be noted that downward compatibility with previous
versions of CP/M is only maintained if the file does not exceed one
megabyte, no file attributes are set, and the file is created by user
0.

6. ED ENHANCEMENTS.

The CP/M standard program editor provides several new facilities in the 2.0 release. Experience has shown that most operators use the relative line numbering feature of ED, and thus the editor has the "v" (Verify Line) option set as an initial value. The operator can, of course, disable line numbering by typing the "-v" command. If you are not familiar with the ED line number mode, you may wish to refer to the Appendix in the ED user's guide, where the "v" command is described.

ED also takes file attributes into account. If the operator attempts to edit a read/only file, the message

** FILE IS READ/ONLY **

appears at the console. The file can be loaded and examined, but cannot be altered in any way. Normally, the operator simply ends the edit session, and uses STAT to change the file attribute to R/W. If the edited file has the "system" attribute set, the message

"SYSTEM" FILE NOT ACCESSIBLE

is displayed at the console, and the edit session is aborted. Again, the STAT program can be used to change the system attribute, if desired.

Finally, the insert mode ("i") command allows CRT line editing functions, as described in Section 2, above.

7. THE XSUB FUNCTION.

An additional utility program is supplied with version 2.0 of
CP/M, called XSUB, which extends the power of the SUBMIT facility to
include line input to programs as well as the console command
processor.   The XSUB command is included as the first line of your
submit file and, when executed, self-relocates directly below the CCP.
All subsequent submit command lines are processed by XSUB, so that
programs which read buffered console input (BDOS function 10) receive
their input directly from the submit file. For example, the file
SAVER.SUB could contain the submit lines:

```
XSUB
DDT
I$1.HEX
R
G0
SAVE 1 $2.COM
```

with a subsequent SUBMIT command:

```
SUBMIT SAVER X Y
```

which substitutes X for $1 and Y for $2 in the command stream.   The
XSUB program loads, followed by DDT which is sent the command lines
"IX.HEX" "R" and "G0" thus returning to the CCP.   The final command
"SAVE 1 Y.COM" is processed by the CCP.

The XSUB program remains in memory, and prints the message

```
(xsub active)
```

on each warm start operation to indicate its presence.   Subsequent
submit command streams do not require the XSUB, unless an intervening
cold start has occurred. Note that XSUB must be loaded after DESPOOL,
if both are to run simultaneously.

8. BDOS INTERFACE CONVENTIONS.

CP/M 2.0 system calls take place in exactly the same manner as earlier versions, with a call to location 0005H, function number in register C, and information address in register pair DE. Single byte values are returned in register A, with double byte values returned in HL (for reasons of compatibility, register A = L and register B = H upon return in all cases). A list of CP/M 2.0 calls is given below, with an asterisk following functions which are either new or revised from version 1.4 to 2.0. Note that a zero value is returned for out-of range function numbers.

0	System Reset	19*	Delete File
1	Console Input	20	Read Sequential
2	Console Output	21	Write Sequential
3	Reader Input	22*	Make File
4	Punch Output	23*	Rename File
5	List Output	24*	Return Login Vector
6*	Direct Console I/O	25	Return Current Disk
7	Get I/O Byte	26	Set DMA Address
8	Set I/O Byte	27	Get Addr(Alloc)
9	Print String	28*	Write Protect Disk
10*	Read Console Buffer	29*	Get Addr(R/O Vector)
11	Get Console Status	30*	Set File Attributes
12*	Return Version Number	31*	Get Addr(Disk Parms)
13	Reset Disk System	32*	Set/Get User Code
14	Select Disk	33*	Read Random
15*	Open File	34*	Write Random
16	Close File	35*	Compute File Size
17*	Search for First	36*	Set Random Record
18*	Search for Next		

(Functions 28, 29, and 32 should be avoided in application programs to maintain upward compatibility with MP/M.) The new or revised functions are described below.

Function 6: Direct Console I/O.

Direct Console I/O is supported under CP/M 2.0 for those applications where it is necessary to avoid the BDOS console I/O operations. Programs which currently perform direct I/O through the BIOS should be changed to use direct I/O under BDOS so that they can be fully supported under future releases of MP/M and CP/M.

Upon entry to function 6, register E either contains hexadecimal FF, denoting a console input request, or register E contains an ASCII character. If the input value is FF, then function 6 returns A = 00 if no character is ready, otherwise A contains the next console input character.

If the input value in E is not FF, then function 6 assumes that E contains a valid ASCII character which is sent to the console.

Function 10: Read Console Buffer.

The console buffer read operation remains unchanged except that console line editing is supported, as described in Section 2. Note also that certain functions which return the carriage to the leftmost position (e.g., ctl-X) do so only to the column position where the prompt ended (previously, the carriage returned to the extreme left margin). This new convention makes operator data input and line correction more legible.

Function 12: Return Version Number.

Function 12 has been redefined to provide information which allows version-independent programming (this was previously the "lift head" function which returned HL=0000 in version 1.4, but performed no operation). The value returned by function 12 is a two-byte value, with H = 00 for the CP/M release (H = 01 for MP/M), and L = 00 for all releases previous to 2.0. CP/M 2.0 returns a hexadecimal 20 in register L, with subsequent version 2 releases in the hexadecimal range 21, 22, through 2F. Using function 12, for example, you can write application programs which provide both sequential and random access functions, with random access disabled when operating under early releases of CP/M.

In the file operations described below, DE addresses a file control block (FCB). Further, all directory operations take place in a reserved area which does not affect write buffers as was the case in version 1.4, with the exception of Search First and Search Next, where compatibility is required.

The File Control Block (FCB) data area consists of a sequence of 33 bytes for sequential access, and a series of 36 bytes in the case that the file is accessed randomly. The default file control block normally located at 005CH can be used for random access files, since bytes 007DH, 007EH, and 007FH are available for this purpose. For notational purposes, the FCB format is shown with the following fields:

```
-----------------------------------------------------------------
|dr|f1|f2|/ /|f8|t1|t2|t3|ex|s1|s2|rc|d0|/ /|dn|cr|r0|r1|r2|
-----------------------------------------------------------------
 00 01 02 ... 08 09 10 11 12 13 14 15 16 ... 31 32 33 34 35
```

where

dr   drive code (0 - 16)
            0 => use default drive for file
            1 => auto disk select drive A,
            2 => auto disk select drive B,
            ...
            16=> auto disk select drive P.

f1...f8   contain the file name in ASCII
            upper case, with high bit = 0

t1,t2,t3   contain the file type in ASCII
            upper case, with high bit = 0
            t1', t2', and t3' denote the
            bit of these positions,
            t1' = 1 => Read/Only file,
            t2' = 1 => SYS file, no DIR list

ex   contains the current extent number,
            normally set to 00 by the user, but
            in range 0 - 31 during file I/O

s1   reserved for internal system use

s2   reserved for internal system use, set
            to zero on call to OPEN, MAKE, SEARCH

rc   record count for extent "ex,"
            takes on values from 0 - 128

d0...dn   filled-in by CP/M, reserved for
            system use

cr   current record to read or write in
            a sequential file operation, normally
            set to zero by user

r0,r1,r2   optional random record number in the
            range 0-65535, with overflow to r2,
            r0,r1 constitute a 16-bit value with
            low byte r0, and high byte r1

     Function 15: Open File.

     The Open File operation is identical to previous definitions,
with the exception that byte s2 is automatically zeroed. Note that
previous versions of CP/M defined this byte as zero, but made no

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checks to assure compliance. Thus, the byte is cleared to ensure upward compatibility with the latest version, where it is required.

Function 17: Search for First.

Search First scans the directory for a match with the file given by the FCB addressed by DE. The value 255 (hexadecimal FF) is returned if the file is not found, otherwise a value of A equal to 0, 1, 2, or 3 is returned indicating the file is present. In the case that the file is found, the current DMA address is filled with the record containing the directory entry, and the relative starting position is A * 32 (i.e., rotate the A register left 5 bits, or ADD A five times). Although not normally required for application programs, the directory information can be extracted from the buffer at this position.

An ASCII question mark (63 decimal, 3F hexadecimal) in any position from f1 through ex matches the corresponding field of any directory entry on the default or auto-selected disk drive. If the dr field contains an ASCII question mark, then the auto disk select function is disabled, the default disk is searched, with the search function returning any matched entry, allocated or free, belonging to any user number. This latter function is not normally used by application programs, but does allow complete flexibility to scan all current directory values. If the dr field is not a question mark, the s2 byte is automatically zeroed.

Function 18: Search for Next.

The Search Next function is similar to the Search First function, except that the directory scan continues from the last matched entry. Similar to function 17, function 18 returns the decimal value 255 in A when no more directory items match.

Function 19: Delete File.

The Delete File function removes files which match the FCB addressed by DE. The filename and type may contain ambiguous references (i.e., question marks in various positions), but the drive select code cannot be ambiguous, as in the Search and Search Next functions.

Function 19 returns a decimal 255 if the reference file or files could not be found, otherwise a value in the range 0 to 3 is returned.

Function 22: Make File.

The Make File operation is identical to previous versions of CP/M, except that byte s2 is zeroed upon entry to the BDOS.

Function 23: Rename File.

The Actions of the file rename functions are the same as previous releases except that the value 255 is returned if the rename function is unsuccessful (the file to rename could not be found), otherwise a value in the range 0 to 3 is returned.

Function 24: Return Login Vector.

The login vector value returned by CP/M 2.0 is a 16-bit value in HL, where the least significant bit of L corresponds to the first drive A, and the high order bit of H corresponds to the sixteenth drive, labelled P. Note that compatibility is maintained with earlier releases, since registers A and L contain the same values upon return.

Function 28: Write Protect Current Disk.

The disk write protect function provides temporary write protection for the currently selected disk. Any attempt to write to the disk, before the next cold or warm start operation produces the message

                    Bdos Err on d: R/O

Function 29: Get R/O Vector.

Function 29 returns a bit vector in register pair HL which indicates drives which have the temporary read/only bit set. Similar to function 24, the least significant bit corresponds to drive A, while the most significant bit corresponds to drive P. The R/O bit is set either by an explicit call to function 28, or by the automatic software mechanisms within CP/M which detect changed disks.

Function 30: Set File Attributes.

The Set File Attributes function allows programmatic manipulation of permanent indicators attached to files. In particular, the R/O and System attributes (t1' and t2' above) can be set or reset. The DE pair addresses an unambiguous file name with the appropriate attributes set or reset. Function 30 searches for a

match, and changes the matched directory entry to contain the selected indicators. Indicators f1' through f4' are not presently used, but may be useful for applications programs, since they are not involved in the matching process during file open and close operations. Indicators f5' through f8' and t3' are reserved for future system expansion.

Function 31: Get Disk Parameter Block Address.

The address of the BIOS resident disk parameter block is returned in HL as a result of this function call. This address can be used for either of two purposes. First, the disk parameter values can be extracted for display and space computation purposes, or transient programs can dynamically change the values of current disk parameters when the disk environment changes, if required. Normally, application programs will not require this facility.

Function 32: Set or Get User Code.

An application program can change or interrogate the currently active user number by calling function 32. If register E = FF hexadecimal, then the value of the current user number is returned in register A, where the value is in the range 0 to 31. If register E is not FF, then the current user number is changed to the value of E (modulo 32).

Function 33: Read Random.

The Read Random function is similar to the sequential file read operation of previous releases, except that the read operation takes place at a particular record number, selected by the 24-bit value constructed from the three byte field following the FCB (byte positions r0 at 33, r1 at 34, and r2 at 35). Note that the sequence of 24 bits is stored with least significant byte first (r0), middle byte next (r1), and high byte last (r2). CP/M release 2.0 does not reference byte r2, except in computing the size of a file (function 35). Byte r2 must be zero, however, since a non-zero value indicates overflow past the end of file.

Thus, in version 2.0, the r0,r1 byte pair is treated as a double-byte, or "word" value, which contains the record to read. This value ranges from 0 to 65535, providing access to any particular record of the 8 megabyte file. In order to process a file using random access, the base extent (extent 0) must first be opened. Although the base extent may or may not contain any allocated data, this ensures that the file is properly recorded in the directory, and is visible in DIR requests. The selected record number is then stored into the random record field (r0,r1), and the BDOS is called to read the record. Upon return from the call, register A either contains an

error code, as listed below, or the value 00 indicating the operation was successful. In the latter case, the current DMA address contains the randomly accessed record. Note that contrary to the sequential read operation, the record number is not advanced. Thus, subsequent random read operations continue to read the same record.

Upon each random read operation, the logical extent and current record values are automatically set. Thus, the file can be sequentially read or written, starting from the current randomly accessed position. Note, however, that in this case, the last randomly read record will be re-read as you switch from random mode to sequential read, and the last record will be re-written as you switch to a sequential write operation. You can, of course, simply advance the random record position following each random read or write to obtain the effect of a sequential I/O operation.

Error codes returned in register A following a random read are listed below.

> 01   reading unwritten data
> 02   (not returned in random mode)
> 03   cannot close current extent
> 04   seek to unwritten extent
> 05   (not returned in read mode)
> 06   seek past physical end of disk

Error code 01 and 04 occur when a random read operation accesses a data block which has not been previously written, or an extent which has not been created, which are equivalent conditions. Error 3 does not normally occur under proper system operation, but can be cleared by simply re-reading, or re-opening extent zero as long as the disk is not physically write protected. Error code 06 occurs whenever byte r2 is non-zero under the current 2.0 release. Normally, non-zero return codes can be treated as missing data, with zero return codes indicating operation complete.

Function 34: Write Random.

The Write Random operation is initiated similar to the Read Random call, except that data is written to the disk from the current DMA address. Further, if the disk extent or data block which is the target of the write has not yet been allocated, the allocation is performed before the write operation continues. As in the Read Random operation, the random record number is not changed as a result of the write. The logical extent number and current record positions of the file control block are set to correspond to the random record which is being written. Again, sequential read or write operations can commence following a random write, with the notation that the currently addressed record is either read or rewritten again as the sequential operation begins. You can also simply advance the random record position following each write to get the effect of a sequential write operation. Note that in particular, reading or writing the last record of an extent in random mode does not cause an automatic extent

switch as it does in sequential mode under either CP/M 1.4 or CP/M 2.0.

The error codes returned by a random write are identical to the random read operation with the addition of error code 05, which indicates that a new extent cannot be created due to directory overflow.

Function 35: Compute File Size.

When computing the size of a file, the DE register pair addresses an FCB in random mode format (bytes r0, r1, and r2 are present). The FCB contains an unambiguous file name which is used in the directory scan. Upon return, the random record bytes contain the "virtual" file size which is, in effect, the record address of the record following the end of the file. if, following a call to function 35, the high record byte r2 is 01, then the file contains the maximum record count 65536 in version 2.0. Otherwise, bytes r0 and r1 constitute a 16-bit value (r0 is the least significant byte, as before) which is the file size.

Data can be appended to the end of an existing file by simply calling function 35 to set the random record position to the end of file, then performing a sequence of random writes starting at the preset record address.

The virtual size of a file corresponds to the physical size when the file is written sequentially. If, instead, the file was created in random mode and "holes" exist in the allocation, then the file may in fact contain fewer records than the size indicates. If, for example, only the last record of an eight megabyte file is written in random mode (i.e., record number 65535), then the virtual size is 65536 records, although only one block of data is actually allocated.

Function 36: Set Random Record.

The Set Random Record function causes the BDOS to automatically produce the random record position from a file which has been read or written sequentially to a particular point. The function can be useful in two ways.

First, it is often necessary to initially read and scan a sequential file to extract the positions of various "key" fields. As each key is encountered, function 36 is called to compute the random record position for the data corresponding to this key. If the data unit size is 128 bytes, the resulting record position is placed into a table with the key for later retrieval. After scanning the entire file and tabularizing the keys and their record numbers, you can move instantly to a particular keyed record by performing a random read using the corresponding random record number which was saved earlier. The scheme is easily generalized when variable record lengths are

involved since the program need only store the buffer-relative byte
position along with the key and record number in order to find the
exact starting position of the keyed data at a later time.

A second use of function 36 occurs when switching from a
sequential read or write over to random read or write. A file is
sequentially accessed to a particular point in the file, function 36
is called which sets the record number, and subsequent random read and
write operations continue from the selected point in the file.

This section is concluded with a rather extensive, but complete
example of random access operation. The program listed below performs
the simple function of reading or writing random records upon command
from the terminal. Given that the program has been created,
assembled, and placed into a file labelled RANDOM.COM, the CCP level
command:

                        RANDOM X.DAT

starts the test program. The program looks for a file by the name
X.DAT (in this particular case) and, if found, proceeds to prompt the
console for input. If not found, the file is created before the
prompt is given. Each prompt takes the form

                        next command?

and is followed by operator input, terminated by a carriage return.
The input commands take the form

                        nW   nR   Q

where n is an integer value in the range 0 to 65535, and W, R, and Q
are simple command characters corresponding to random write, random
read, and quit processing, respectively. If the W command is issued,
the RANDOM program issues the prompt

                        type data:

The operator then responds by typing up to 127 characters, followed by
a carriage return. RANDOM then writes the character string into the
X.DAT file at record n. If the R command is issued, RANDOM reads
record number n and displays the string value at the console. If the
Q command is issued, the X.DAT file is closed, and the program returns
to the console command processor. In the interest of brevity (ok, so
the program's not so brief), the only error message is

                        error, try again

The program begins with an initialization section where the
input file is opened or created, followed by a continuous loop at the
label "ready" where the individual commands are interpreted. The
default file control block at 005CH and the default buffer at 0080H
are used in all disk operations. The utility subroutines then follow,

which contain the principal input line processor, called "readc."
This particular program shows the elements of random access
processing, and can be used as the basis for further program
development.

```
                ;*********************************************************
                ;*                                                     *
                ;* sample random access program for cp/m 2.0           *
                ;*                                                     *
                ;*********************************************************
0100                    org     100h        ;base of tpa
                ;
0000 =          reboot  equ     0000h       ;system reboot
0005 =          bdos    equ     0005h       ;bdos entry point
                ;
0001 =          coninp  equ     1           ;console input function
0002 =          conout  equ     2           ;console output function
0009 =          pstring equ     9           ;print string until '$'
000a =          rstring equ     10          ;read console buffer
000c =          version equ     12          ;return version number
000f =          openf   equ     15          ;file open function
0010 =          closef  equ     16          ;close function
0016 =          makef   equ     22          ;make file function
0021 =          readr   equ     33          ;read random
0022 =          writer  equ     34          ;write random
                ;
005c =          fcb     equ     005ch       ;default file control block
007d =          ranrec  equ     fcb+33      ;random record position
007f =          ranovf  equ     fcb+35      ;high order (overflow) byte
0080 =          buff    equ     0080h       ;buffer address
                ;
000d =          cr      equ     0dh         ;carriage return
000a =          lf      equ     0ah         ;line feed
                ;
                ;*********************************************************
                ;*                                                     *
                ;* load SP, set-up file for random access              *
                ;*                                                     *
                ;*********************************************************
0100 31bc0              lxi     sp,stack
                ;
                ;               version 2.0?
0103 0e0c              mvi     c,version
0105 cd0500            call    bdos
0108 fe20             cpi     20h         ;version 2.0 or better?
010a d21600           jnc     versok
                ;       bad version, message and go back
010d 111b0             lxi     d,badver
0110 cdda0            call    print
0113 c3000            jmp     reboot
                ;
                versok:
                ;               correct version for random access
```

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```
0116 0e0f            mvi     c,openf ;open default fcb
0118 115c0           lxi     d,fcb
011b cd050           call    bdos
011e 3c              inr     a       ;err 255 becomes zero
011f c2370           jnz     ready
            ;
            ;               cannot open file, so create it
0122 0e16            mvi     c,makef
0124 115c0           lxi     d,fcb
0127 cd050           call    bdos
012a 3c              inr     a       ;err 255 becomes zero
012b c2370           jnz     ready
            ;
            ;               cannot create file, directory full
012e 113a0           lxi     d,nospace
0131 cdda0           call    print
0134 c3000           jmp     reboot  ;back to ccp
            ;
            ;********************************************************
            ;*                                                      *
            ;*   loop back to "ready" after each command            *
            ;*                                                      *
            ;********************************************************
            ;
            ready:
            ;               file is ready for processing
            ;
0137 cde50           call    readcom ;read next command
013a 227d0           snld    ranrec  ;store input record#
013d 217f0           lxi     h,ranovf
0140 3600            mvi     m,0     ;clear high byte if set
0142 fe51            cpi     'Q'     ;quit?
0144 c2560           jnz     notq
            ;
            ;               quit processing, close file
0147 0e10            mvi     c,closef
0149 115c0           lxi     d,fcb
014c cd050           call    bdos
014f 3c              inr     a       ;err 255 becomes 0
0150 cab90           jz      error   ;error message, retry
0153 c3000           jmp     reboot  ;back to ccp
            ;
            ;********************************************************
            ;*                                                      *
            ;* end of quit command, process write                   *
            ;*                                                      *
            ;********************************************************
            notq:
            ;               not the quit command, random write?
0156 fe57            cpi     'W'
0158 c2890           jnz     notw
            ;
            ;               this is a random write, fill buffer until cr
015b 114d0           lxi     d,datmsg
015e cdda0           call    print   ;data prompt
```

```
0161 0e7f              mvi     c,127    ;up to 127 characters
0163 21800             lxi     h,buff   ;destination
            rloop:     ;read next character to buff
0166 c5                push    b        ;save counter
0167 e5                push    h        ;next destination
0168 cdc20             call    getchr   ;character to a
016b e1                pop     h        ;restore counter
016c c1                pop     b        ;restore next to fill
016d fe0d              cpi     cr       ;end of line?
016f ca780             jz      erloop
            ;          not end, store character
0172 77                mov     m,a
0173 23                inx     h        ;next to fill
0174 0d                dcr     c        ;counter goes down
0175 c2660             jnz     rloop    ;end of buffer?
            erloop:
            ;          end of read loop, store 00
0178 3600              mvi     m,0
            ;
            ;          write the record to selected record number
017a 0e22              mvi     c,writer
017c 115c0             lxi     d,fcb
017f cd050             call    bdos
0182 b7                ora     a        ;error code zero?
0183 c2b90             jnz     error    ;message if not
0186 c3370             jmp     ready    ;for another record
            ;
            ;****************************************************
            ;*                                                *
            ;* end of write command, process read             *
            ;*                                                *
            ;****************************************************
            notw:
            ;          not a write command, read record?
0189 fe52              cpi     'R'
018b c2b90             jnz     error    ;skip if not
            ;
            ;          read random record
018e 0e21              mvi     c,readr
0190 115c0             lxi     d,fcb
0193 cd050             call    bdos
0196 b7                ora     a        ;return code 00?
0197 c2b90             jnz     error
            ;
            ;          read was successful, write to console
019a cdcf0             call    crlf     ;new line
019d 0e80              mvi     c,128    ;max 128 characters
019f 21800             lxi     h,buff   ;next to get
            wloop:
01a2 7e                mov     a,m      ;next character
01a3 23                inx     h        ;next to get
01a4 e67f              ani     7fh      ;mask parity
01a6 ca370             jz      ready    ;for another command if 00
01a9 c5                push    b        ;save counter
01aa e5                push    h        ;save next to get
```

```
01ab  fe20               cpi                  ;graphic?
01ad  d4c80              cnc      putchr      ;skip output if not
01b0  el                 pop      h
01b1  cl                 pop      b
01b2  0d                 dcr      c           ;count=count-1
01b3  c2a20              jnz      wloop
01b6  c3370              jmp      ready
                 ;
                 ;*********************************************************
                 ;*                                                      *
                 ;* end of read command, all errors end-up here          *
                 ;*                                                      *
                 ;*********************************************************
                 ;
                 error:
01b9  11590              lxi      d,errmsg
01bc  cdda0              call     print
01bf  c3370              jmp      ready
                 ;
                 ;*********************************************************
                 ;*                                                      *
                 ;* utility subroutines for console i/o                  *
                 ;*                                                      *
                 ;*********************************************************
                 getchr:
                                  ;read next console character to a
01c2  0e01               mvi      c,coninp
01c4  cd050              call     bdos
01c7  c9                 ret
                 ;
                 putchr:
                                  ;write character from a to console
01c8  0e02               mvi      c,conout
01ca  5f                 mov      e,a         ;character to send
01cb  cd050              call     bdos        ;send character
01ce  c9                 ret
                 ;
                 crlf:
                                  ;send carriage return line feed
01cf  3e0d               mvi      a,cr        ;carriage return
01d1  cdc80              call     putchr
01d4  3e0a               mvi      a,lf        ;line feed
01d6  cdc80              call     putchr
01d9  c9                 ret
                 ;
                 print:
                                  ;print the buffer addressed by de until $
01da  d5                 push     d
01db  cdcf0              call     crlf
01de  dl                 pop      d           ;new line
01df  0e09               mvi      c,pstring
01el  cd050              call     bdos        ;print the string
01e4  c9                 ret
                 ;
                 readcom:
```

```
                            ;read the next command line to the conbuf
    01e5 116b0      lxi     d,prompt
    01e8 cdda0      call    print    ;command?
    01eb 0e0a       mvi     c,rstring
    01ed 117a0      lxi     d,conbuf
    01f0 cd050      call    bdos     ;read command line
             ;      command line is present, scan it
    01f3 21000      lxi     h,0      ;start with 0000
    01f6 117c0      lxi     d,conlin ;command line
    01f9 1a   readc: ldax   d        ;next command character
    01fa 13         inx     d        ;to next command position
    01fb b7         ora     a        ;cannot be end of command
    01fc c8         rz
             ;      not zero, numeric?
    01fd d630       sui     '0'
    01ff fe0a       cpi     10       ;carry if numeric
    0201 d2130      jnc     endrd
             ;      add-in next digit
    0204 29         dad     h        ;*2
    0205 4d         mov     c,l
    0206 44         mov     b,h      ;bc = value * 2
    0207 29         dad     h        ;*4
    0208 29         dad     h        ;*8
    0209 09         dad     b        ;*2 + *8 = *10
    020a 85         add     l        ;+digit
    020b 6f         mov     l,a
    020c d2f90      jnc     readc    ;for another char
    020f 24         inr     h        ;overflow
    0210 c3f90      jmp     readc    ;for another char
          endrd:
             ;      end of read, restore value in a
    0213 c630       adi     '0'      ;command
    0215 fe61       cpi     'a'      ;translate case?
    0217 d8         rc
             ;      lower case, mask lower case bits
    0218 e65f       ani     101$1111b
    021a c9         ret
          ;
          ;************************************************************
          ;*                                                          *
          ;* string data area for console messages                   *
          ;*                                                          *
          ;************************************************************
          badver:
    021b 536f79     db      'sorry, you need cp/m version 2$'
          nospace:
    023a 4e6f29     db      'no directory space$'
          datmsg:
    024d 547970     db      'type data: $'
          errmsg:
    0259 457272     db      'error, try again.$'
          prompt:
    026b 4e6570     db      'next command? $'
          ;
```

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```
        ;****************************************************
        ;*                                                  *
        ;*  fixed and variable data area                    *
        ;*                                                  *
        ;****************************************************
027a 21 conbuf: db      conlen   ;length of console buffer
027b    consiz: ds      1        ;resulting size after read
027c    conlin: ds      32       ;length 32 buffer
0021 =  conlen  equ     $-consiz
        ;
029c            ds      32       ;16 level stack
        stack:
02bc            end
```