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CP/M 2.2 ALTERATION GUIDE

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1. INTRODUCTION

The standard CP/M system assumes operation on an Intel MDS-800 microcomputer development system, but is designed so that the user can alter a specific set of subroutines which define the hardware operating environment. In this way, the user can produce a diskette which operates with any IBM-3741 format compatible drive controller and other peripheral devices.

Although standard CP/M 2.0 is configured for single density floppy disks, field-alteration features allow adaptation to a wide variety of disk subsystems from single drive minidisks through high-capacity "hard disk" systems. In order to simplify the following adaptation process, we assume that CP/M 2.0 will first be configured for single density floppy disks where minimal editing and debugging tools are available. If an earlier version of CP/M is available, the customizing process is eased considerably. In this latter case, you may wish to briefly review the system generation process, and skip to later sections which discuss system alteration for non-standard disk systems.

In order to achieve device independence, CP/M is separated into three distinct modules:

- BIOS - basic I/O system which is environment dependent
- BDOS - basic disk operating system which is not dependent upon the hardware configuration
- CCP - the console command processor which uses the BDOS

Of these modules, only the BIOS is dependent upon the particular hardware. That is, the user can "patch" the distribution version of CP/M to provide a new BIOS which provides a customized interface between the remaining CP/M modules and the user's own hardware system. The purpose of this document is to provide a step-by-step procedure for patching your new BIOS into CP/M.

If CP/M is being tailored to your computer system for the first time, the new BIOS requires some relatively simple software development and testing. The standard BIOS is listed in Appendix B, and can be used as a model for the customized package. A skeletal version of the BIOS is given in Appendix C which can serve as the basis for a modified BIOS. In addition to the BIOS, the user must write a simple memory loader, called GETSYS, which brings the operating system into memory. In order to patch the new BIOS into CP/M, the user must write the reverse of GETSYS, called PUTSYS, which places an altered version of CP/M back onto the diskette. PUTSYS can be derived from GETSYS by changing the disk read commands into disk write commands. Sample skeletal GETSYS and PUTSYS programs are described in Section 3, and listed in Appendix D. In order to make the CP/M system work automatically, the user must also supply a cold start loader, similar to the one provided with CP/M (listed in Appendices A and B). A skeletal form of a cold start loader is given in Appendix E which can serve as a model for your loader.

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2. FIRST LEVEL SYSTEM REGENERATION

The procedure to follow to patch the CP/M system is given below in several steps. Address references in each step are shown with a following "H" which denotes the hexadecimal radix, and are given for a 20K CP/M system. For larger CP/M systems, add a "bias" to each address which is shown with a "+b" following it, where b is equal to the memory size - 20K. Values for b in various standard memory sizes are

24K:	b = 24K - 20K = 4K = 1000H
32K:	b = 32K - 20K = 12K = 3000H
40K:	b = 40K - 20K = 20K = 5000H
48K:	b = 48K - 20K = 28K = 7000H
56K:	b = 56K - 20K = 36K = 9000H
62K:	b = 62K - 20K = 42K = A800H
64K:	b = 64K - 20K = 44K = B000H

Note: The standard distribution version of CP/M is set for operation within a 20K memory system. Therefore, you must first bring up the 20K CP/M system, and then configure it for your actual memory size (see Second Level System Generation).

(1) Review Section 4 and write a GETSYS program which reads the first two tracks of a diskette into memory. The data from the diskette must begin at location 3380H. Code GETSYS so that it starts at location 100H (base of the TPA), as shown in the first part of Appendix d.

(2) Test the GETSYS program by reading a blank diskette into memory, and check to see that the data has been read properly, and that the diskette has not been altered in any way by the GETSYS program.

(3) Run the GETSYS program using an initialized CP/M diskette to see if GETSYS loads CP/M starting at 3380H (the operating system actually starts 128 bytes later at 3400H).

(4) Review Section 4 and write the PUTSYS program which writes memory starting at 3380H back onto the first two tracks of the diskette. The PUTSYS program should be located at 200H, as shown in the second part of Appendix D.

(5) Test the PUTSYS program using a blank uninitialized diskette by writing a portion of memory to the first two tracks; clear memory and read it back using GETSYS. Test PUTSYS completely, since this program will be used to alter CP/M on disk.

(6) Study Sections 5, 6, and 7, along with the distribution version of the BIOS given in Appendix B, and write a simple version which performs a similar function for the customized environment. Use the program given in Appendix C as a model. Call this new BIOS by the name CBIOS (customized BIOS). Implement only the primitive disk operations on a single drive, and simple console input/output functions in this phase.

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(7) Test CBIOS completely to ensure that it properly performs console character I/O and disk reads and writes. Be especially careful to ensure that no disk write operations occur accidentally during read operations, and check that the proper track and sectors are addressed on all reads and writes. Failure to make these checks may cause destruction of the initialized CP/M system after it is patched.

(8) Referring to Figure 1 in Section 5, note that the BIOS is placed between locations 4A00H and 4FFFH. Read the CP/M system using GETSYS and replace the BIOS segment by the new CBIOS developed in step (6) and tested in step (7). This replacement is done in the memory of the machine, and will be placed on the diskette in the next step.

(9) Use PUTSYS to place the patched memory image of CP/M onto the first two tracks of a blank diskette for testing.

(10) Use GETSYS to bring the copied memory image from the test diskette back into memory at 3380H, and check to ensure that it has loaded back properly (clear memory, if possible, before the load). Upon successful load, branch to the cold start code at location 4A00H. The cold start routine will initialize page zero, then jump to the CCP at location 3400H which will call the BDOS, which will call the CBIOS. The CBIOS will be asked by the CCP to read sixteen sectors on track 2, and if successful, CP/M will type "A>", the system prompt.

When you make it this far, you are almost on the air. If you have trouble, use whatever debug facilities you have available to trace and breakpoint your CBIOS.

(11) Upon completion of step (10), CP/M has prompted the console for a command input. Test the disk write operation by typing

```
SAVE 1 X.COM
```

(recall that all commands must be followed by a carriage return).

CP/M should respond with another prompt (after several disk accesses):

```
A>
```

If it does not, debug your disk write functions and retry.

(12) Then test the directory command by typing

```
DIR
```

CP/M should respond with

```
A: X      COM
```

(13) Test the erase command by typing

```
ERA X.COM
```

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CP/M should respond with the A prompt. When you make it this far, you should have an operational system which will only require a bootstrap loader to function completely.

(14) Write a bootstrap loader which is similar to GETSYS, and place it on track 0, sector 1 using PUTSYS (again using the test diskette, not the distribution diskette). See Sections 5 and 8 for more information on the bootstrap operation.

(15) Retest the new test diskette with the bootstrap loader installed by executing steps (11), (12), and (13). Upon completion of these tests, type a control-C (control and C keys simultaneously). The system should then execute a "warm start" which reboots the system, and types the A prompt.

(16) At this point, you probably have a good version of your customized CP/M system on your test diskette. Use GETSYS to load CP/M from your test diskette. Remove the test diskette, place the distribution diskette (or a legal copy) into the drive, and use PUTSYS to replace the distribution version by your customized version. Do not make this replacement if you are unsure of your patch since this step destroys the system which was sent to you from Digital Research.

(17) Load your modified CP/M system and test it by typing

DIR

CP/M should respond with a list of files which are provided on the initialized diskette. One such file should be the memory image for the debugger, called DDT.COM.

NOTE: from now on, it is important that you always reboot the CP/M system (ctl-C is sufficient) when the diskette is removed and replaced by another diskette, unless the new diskette is to be read only.

(18) Load and test the debugger by typing

DDT

(see the document "CP/M Dynamic Debugging Tool (DDT)" for operating procedures. You should take the time to become familiar with DDT, it will be your best friend in later steps.

(19) Before making further CBIOS modifications, practice using the editor (see the ED user's guide), and assembler (see the ASM user's guide). Then recode and test the GETSYS, PUTSYS, and CBIOS programs using ED, ASM, and DDT. Code and test a COPY program which does a sector-to-sector copy from one diskette to another to obtain back-up copies of the original diskette (NOTE: read your CP/M Licensing Agreement; it specifies your legal responsibilities when copying the CP/M system). Place the copyright notice

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on each copy which is made with your COPY program.

(20) Modify your CBIOS to include the extra functions for punches, readers, signon messages, and so-forth, and add the facilities for a additional disk drives, if desired. You can make these changes with the GETSYS and PUTSYS programs which you have developed, or you can refer to the following section, which outlines CP/M facilities which will aid you in the regeneration process.

You now have a good copy of the customized CP/M system. Note that although the CBIOS portion of CP/M which you have developed belongs to you, the modified version of CP/M which you have created can be copied for your use only (again, read your Licensing Agreement), and cannot be legally copied for anyone else's use.

It should be noted that your system remains file-compatible with all other CP/M systems, (assuming media compatibility, of course) which allows transfer of non-proprietary software between users of CP/M.

3. SECOND LEVEL SYSTEM GENERATION

Now that you have the CP/M system running, you will want to configure CP/M for your memory size. In general, you will first get a memory image of CP/M with the "MOVCPM" program (system relocater) and place this memory image into a named disk file. The disk file can then be loaded, examined, patched, and replaced using the debugger, and system generation program. For further details on the operation of these programs, see the "Guide to CP/M Features and Facilities" manual.

Your CBIOS and BOOT can be modified using ED, and assembled using ASM, producing files called CBIOS.HEX and BOOT.HEX, which contain the machine code for CBIOS and BOOT in Intel hex format.

To get the memory image of CP/M into the TPA configured for the desired memory size, give the command:

```
MOVCPM xx *
```

where "xx" is the memory size in decimal K bytes (e.g., 32 for 32K). The response will be:

```
CONSTRUCTING xxK CP/M VERS 2.0  
READY FOR "SYSGEN" OR  
"SAVE 34 CPMxx.COM"
```

At this point, an image of a CP/M in the TPA configured for the requested memory size. The memory image is at location 0900H through 227FH. (i.e., The BOOT is at 0900H, the CCP is at 980H, the BDOS starts at 1180H, and the BIOS is at 1F80H.) Note that the memory image has the standard MDS-800 BIOS and BOOT on it. It is now necessary to save the memory image in a file so that you can patch your CBIOS and CBOOT into it:

```
SAVE 34 CPMxx.COM
```

The memory image created by the "MOVCPM" program is offset by a negative bias so that it loads into the free area of the TPA, and thus does not interfere with the operation of CP/M in higher memory. This memory image can be subsequently loaded under DDT and examined or changed in preparation for a new generation of the system. DDT is loaded with the memory image by typing:

```
DDT CPMxx.COM
```

Load DDT, then read the CP image

DDT should respond with

```
NEXT PC  
2300 0100  
-
```

(The DDT prompt)

You can then use the display and disassembly commands to examine

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portions of the memory image between 900H and 227FH. Note, however, that to find any particular address within the memory image, you must apply the negative bias to the CP/M address to find the actual address. Track 00, sector 01 is loaded to location 900H (you should find the cold start loader at 900H to 97FH), track 00, sector 02 is loaded into 980H (this is the base of the CCP), and so-forth through the entire CP/M system load. In a 20K system, for example, the CCP resides at the CP/M address 3400H, but is placed into memory at 980H by the SYSGEN program. Thus, the negative bias, denoted by n, satisfies

$$3400H + n = 980H, \text{ or } n = 980H - 3400H$$

Assuming two's complement arithmetic, $n = D580H$, which can be checked by

$$3400H + D580H = 10980H = 0980H \text{ (ignoring high-order overflow).}$$

Note that for larger systems, n satisfies

$$\begin{aligned} (3400H+b) + n &= 980H, \text{ or} \\ n &= 980H - (3400H + b), \text{ or} \\ n &= D580H - b. \end{aligned}$$

The value of n for common CP/M systems is given below

memory size	bias b	negative offset n
20K	0000H	D580H - 0000H = D580H
24K	1000H	D580H - 1000H = C580H
32K	3000H	D580H - 3000H = A580H
40K	5000H	D580H - 5000H = 8580H
48K	7000H	D580H - 7000H = 6580H
56K	9000H	D580H - 9000H = 4580H
62K	A800H	D580H - A800H = 2D80H
64K	B000H	D580H - B000H = 2580H

Assume, for example, that you want to locate the address x within the memory image loaded under DDT in a 20K system. First type

Hx,n Hexadecimal sum and difference

and DDT will respond with the value of x+n (sum) and x-n (difference). The first number printed by DDT will be the actual memory address in the image where the data or code will be found. The input

H3400,D580

for example, will produce 980H as the sum, which is where the CCP is located in the memory image under DDT.

Use the L command to disassemble portions the BIOS located at $(4A00H+b)-n$ which, when you use the H command, produces an actual address of 1F80H. The disassembly command would thus be

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SYSGEN	Start the SYSGEN program
SYSGEN VERSION 2.0	Sign-on message from SYSGEN
SOURCE DRIVE NAME (OR RETURN TO SKIP)	Respond with a carriage return to skip the CP/M read operation since the system is already in memory.
DESTINATION DRIVE NAME (OR RETURN TO REBOOT)	Respond with "B" to write the new system to the diskette in drive B.
DESTINATION ON B, THEN TYPE RETURN	Place a scratch diskette in drive B, then type return.
FUNCTION COMPLETE	
DESTINATION DRIVE NAME (OR RETURN TO REBOOT)	

Place the scratch diskette in your drive A, and then perform a coldstart to bring up the new CP/M system you have configured.

Test the new CP/M system, and place the Digital Research copyright notice on the diskette, as specified in your Licensing Agreement:

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4. SAMPLE GETSYS AND PUTSYS PROGRAMS

The following program provides a framework for the GETSYS and PUTSYS programs referenced in Section 2. The READSEC and WRITESEC subroutines must be inserted by the user to read and write the specific sectors.

```

; GETSYS PROGRAM - READ TRACKS 0 AND 1 TO MEMORY AT 3380H
; REGISTER USE
; A (SCRATCH REGISTER)
; B TRACK COUNT (0, 1)
; C SECTOR COUNT (1,2,...,26)
; DE (SCRATCH REGISTER PAIR)
; HL LOAD ADDRESS
; SP SET TO STACK ADDRESS
;
START: LXI SP,3380H ;SET STACK POINTER TO SCRATCH AREA
       LXI H, 3380H ;SET BASE LOAD ADDRESS
       MVI B, 0 ;START WITH TRACK 0
RDTRK: ;READ NEXT TRACK (INITIALLY 0)
       MVI C,1 ;READ STARTING WITH SECTOR 1
RDSEC: ;READ NEXT SECTOR
       CALL READSEC ;USER-SUPPLIED SUBROUTINE
       LXI D,128 ;MOVE LOAD ADDRESS TO NEXT 1/2 PAGE
       DAD D ;HL = HL + 128
       INR C ;SECTOR = SECTOR + 1
       MOV A,C ;CHECK FOR END OF TRACK
       CPI 27
       JC RDSEC ;CARRY GENERATED IF SECTOR < 27
;
; ARRIVE HERE AT END OF TRACK, MOVE TO NEXT TRACK
       INR B
       MOV A,B ;TEST FOR LAST TRACK
       CPI 2
       JC RDTRK ;CARRY GENERATED IF TRACK < 2
;
; ARRIVE HERE AT END OF LOAD, HALT FOR NOW
       HLT
;
; USER-SUPPLIED SUBROUTINE TO READ THE DISK
READSEC:
; ENTER WITH TRACK NUMBER IN REGISTER B,
; SECTOR NUMBER IN REGISTER C, AND
; ADDRESS TO FILL IN HL
;
       PUSH B ;SAVE B AND C REGISTERS
       PUSH H ;SAVE HL REGISTERS
       .....
       perform disk read at this point, branch to
       label START if an error occurs
       .....
       POP H ;RECOVER HL
       POP B ;RECOVER B AND C REGISTERS
       RET ;BACK TO MAIN PROGRAM
END START

```

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Note that this program is assembled and listed in Appendix C for reference purposes, with an assumed origin of 100H. The hexadecimal operation codes which are listed on the left may be useful if the program has to be entered through your machine's front panel switches.

The PUTSYS program can be constructed from GETSYS by changing only a few operations in the GETSYS program given above, as shown in Appendix D. The register pair HL become the dump address (next address to write), and operations upon these registers do not change within the program. The READSEC subroutine is replaced by a WRITESEC subroutine which performs the opposite function: data from address HL is written to the track given by register B and sector given by register C. It is often useful to combine GETSYS and PUTSYS into a single program during the test and development phase, as shown in the Appendix.

5. DISKETTE ORGANIZATION

The sector allocation for the standard distribution version of CP/M is given here for reference purposes. The first sector (see table on the following page) contains an optional software boot section. Disk controllers are often set up to bring track 0, sector 1 into memory at a specific location (often location 0000H). The program in this sector, called BOOT, has the responsibility of bringing the remaining sectors into memory starting at location 3400H+b. If your controller does not have a built-in sector load, you can ignore the program in track 0, sector 1, and begin the load from track 0 sector 2 to location 3400H+b.

As an example, the Intel MDS-800 hardware cold start loader brings track 0, sector 1 into absolute address 3000H. Upon loading this sector, control transfers to location 3000H, where the bootstrap operation commences by loading the remainder of tracks 0, and all of track 1 into memory, starting at 3400H+b. The user should note that this bootstrap loader is of little use in a non-MDS environment, although it is useful to examine it since some of the boot actions will have to be duplicated in your cold start loader.

Track#	Sector#	Page#	Memory Address	CP/M Module name
00	01		(boot address)	Cold Start Loader
00	02	00	3400H+b	CCP
"	03	"	3480H+b	"
"	04	01	3500H+b	"
"	05	"	3580H+b	"
"	06	02	3600H+b	"
"	07	"	3680H+b	"
"	08	03	3700H+b	"
"	09	"	3780H+b	"
"	10	04	3800H+b	"
"	11	"	3880H+b	"
"	12	05	3900H+b	"
"	13	"	3980H+b	"
"	14	06	3A00H+b	"
"	15	"	3A80H+b	"
"	16	07	3B00H+b	"
00	17	"	3B80H+b	CCP
00	18	08	3C00H+b	BDOS
"	19	"	3C80H+b	"
"	20	09	3D00H+b	"
"	21	"	3D80H+b	"
"	22	10	3E00H+b	"
"	23	"	3E80H+b	"
"	24	11	3F00H+b	"
"	25	"	3F80H+b	"
"	26	12	4000H+b	"
01	01	"	4080H+b	"
"	02	13	4100H+b	"
"	03	"	4180H+b	"
"	04	14	4200H+b	"
"	05	"	4280H+b	"
"	06	15	4300H+b	"
"	07	"	4380H+b	"
"	08	16	4400H+b	"
"	09	"	4480H+b	"
"	10	17	4500H+b	"
"	11	"	4580H+b	"
"	12	18	4600H+b	"
"	13	"	4680H+b	"
"	14	19	4700H+b	"
"	15	"	4780H+b	"
"	16	20	4800H+b	"
"	17	"	4880H+b	"
"	18	21	4900H+b	"
01	19	"	4980H+b	BDOS
01	20	22	4A00H+b	BIOS
"	21	"	4A80H+b	"
"	23	23	4B00H+b	"
"	24	"	4B80H+b	"
"	25	24	4C00H+b	"
01	26	"	4C80H+b	BIOS
02-76	01-26			(directory and data)

6. THE BIOS ENTRY POINTS

The entry points into the BIOS from the cold start loader and BDOS are detailed below. Entry to the BIOS is through a "jump vector" located at 4A00H+b, as shown below (see Appendices B and C, as well). The jump vector is a sequence of 17 jump instructions which send program control to the individual BIOS subroutines. The BIOS subroutines may be empty for certain functions (i.e., they may contain a single RET operation) during regeneration of CP/M, but the entries must be present in the jump vector.

The jump vector at 4A00H+b takes the form shown below, where the individual jump addresses are given to the left:

4A00H+b	JMP BOOT	; ARRIVE HERE FROM COLD START LOAD
4A03H+b	JMP WBOOT	; ARRIVE HERE FOR WARM START
4A06H+b	JMP CONST	; CHECK FOR CONSOLE CHAR READY
4A09H+b	JMP CONIN	; READ CONSOLE CHARACTER IN
4A0CH+b	JMP CONOUT	; WRITE CONSOLE CHARACTER OUT
4A0FH+b	JMP LIST	; WRITE LISTING CHARACTER OUT
4A12H+b	JMP PUNCH	; WRITE CHARACTER TO PUNCH DEVICE
4A15H+b	JMP READER	; READ READER DEVICE
4A18H+b	JMP HOME	; MOVE TO TRACK 00 ON SELECTED DISK
4A1BH+b	JMP SELDSK	; SELECT DISK DRIVE
4A1EH+b	JMP SETTRK	; SET TRACK NUMBER
4A21H+b	JMP SETSEC	; SET SECTOR NUMBER
4A24H+b	JMP SETDMA	; SET DMA ADDRESS
4A27H+b	JMP READ	; READ SELECTED SECTOR
4A2AH+b	JMP WRITE	; WRITE SELECTED SECTOR
4A2DH+b	JMP LISTST	; RETURN LIST STATUS
4A30H+b	JMP SECTAN	; SECTOR TRANSLATE SUBROUTINE

Each jump address corresponds to a particular subroutine which performs the specific function, as outlined below. There are three major divisions in the jump table: the system (re)initialization which results from calls on BOOT and WBOOT, simple character I/O performed by calls on CONST, CONIN, CONOUT, LIST, PUNCH, READER, and LISTST, and diskette I/O performed by calls on HOME, SELDSK, SETTRK, SETSEC, SETDMA, READ, WRITE, and SECTAN.

All simple character I/O operations are assumed to be performed in ASCII, upper and lower case, with high order (parity bit) set to zero. An end-of-file condition for an input device is given by an ASCII control-z (1AH). Peripheral devices are seen by CP/M as "logical" devices, and are assigned to physical devices within the BIOS.

In order to operate, the BDOS needs only the CONST, CONIN, and CONOUT subroutines (LIST, PUNCH, and READER may be used by PIP, but not the BDOS). Further, the LISTST entry is used currently only by DESPOOL, and thus, the initial version of CBIOS may have empty subroutines for the remaining ASCII devices.

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The characteristics of each device are

- CONSOLE The principal interactive console which communicates with the operator, accessed through CONST, CONIN, and CONOUT. Typically, the CONSOLE is a device such as a CRT or Teletype.
- LIST The principal listing device, if it exists on your system, which is usually a hard-copy device, such as a printer or Teletype.
- PUNCH The principal tape punching device, if it exists, which is normally a high-speed paper tape punch or Teletype.
- READER The principal tape reading device, such as a simple optical reader or Teletype.

Note that a single peripheral can be assigned as the LIST, PUNCH, and READER device simultaneously. If no peripheral device is assigned as the LIST, PUNCH, or READER device, the CBIOS created by the user may give an appropriate error message so that the system does not "hang" if the device is accessed by PIP or some other user program. Alternately, the PUNCH and LIST routines can just simply return, and the READER routine can return with a LAH (ctl-Z) in reg A to indicate immediate end-of-file.

For added flexibility, the user can optionally implement the "IOBYTE" function which allows reassignment of physical and logical devices. The IOBYTE function creates a mapping of logical to physical devices which can be altered during CP/M processing (see the STAT command). The definition of the IOBYTE function corresponds to the Intel standard as follows: a single location in memory (currently location 0003H) is maintained, called IOBYTE, which defines the logical to physical device mapping which is in effect at a particular time. The mapping is performed by splitting the IOBYTE into four distinct fields of two bits each, called the CONSOLE, READER, PUNCH, and LIST fields, as shown below:

	most significant			least significant
IOBYTE AT 0003H	LIST	PUNCH	READER	CONSOLE
	bits 6,7	bits 4,5	bits 2,3	bits 0,1

The value in each field can be in the range 0-3, defining the assigned source or destination of each logical device. The values which can be assigned to each field are given below

CONSOLE field (bits 0,1)

- 0 - console is assigned to the console printer device (TTY:)
- 1 - console is assigned to the CRT device (CRT:)
- 2 - batch mode: use the READER as the CONSOLE input, and the LIST device as the CONSOLE output (BAT:)
- 3 - user defined console device (UC1:)

READER field (bits 2,3)

- 0 - READER is the Teletype device (TTY:)
- 1 - READER is the high-speed reader device (RDR:)
- 2 - user defined reader # 1 (UR1:)
- 3 - user defined reader # 2 (UR2:)

PUNCH field (bits 4,5)

- 0 - PUNCH is the Teletype device (TTY:)
- 1 - PUNCH is the high speed punch device (PUN:)
- 2 - user defined punch # 1 (UP1:)
- 3 - user defined punch # 2 (UP2:)

LIST field (bits 6,7)

- 0 - LIST is the Teletype device (TTY:)
- 1 - LIST is the CRT device (CRT:)
- 2 - LIST is the line printer device (LPT:)
- 3 - user defined list device (UL1:)

Note again that the implementation of the IOBYTE is optional, and affects only the organization of your CBIOS. No CP/M systems use the IOBYTE (although they tolerate the existence of the IOBYTE at location 0003H), except for PIP which allows access to the physical devices, and STAT which allows logical-physical assignments to be made and/or displayed (for more information, see the "CP/M Features and Facilities Guide"). In any case, the IOBYTE implementation should be omitted until your basic CBIOS is fully implemented and tested; then add the IOBYTE to increase your facilities.

Disk I/O is always performed through a sequence of calls on the various disk access subroutines which set up the disk number to access, the track and sector on a particular disk, and the direct memory access (DMA) address involved in the I/O operation. After all these parameters have been set up, a call is made to the READ or WRITE function to perform the actual I/O operation. Note that there is often a single call to SELDSK to select a disk drive, followed by a number of read or write operations to the selected disk before selecting another drive for subsequent operations. Similarly, there may be a single call to set the DMA address, followed by several calls which read or write from the selected DMA address before the DMA address is changed. The track and sector subroutines are always called before the READ or WRITE operations are performed.

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Note that the READ and WRITE routines should perform several retries (10 is standard) before reporting the error condition to the BDOS. If the error condition is returned to the BDOS, it will report the error to the user. The HOME subroutine may or may not actually perform the track 00 seek, depending upon your controller characteristics; the important point is that track 00 has been selected for the next operation, and is often treated in exactly the same manner as SETTRK with a parameter of 00.

The exact responsibilities of each entry point subroutine are given below:

BOOT The BOOT entry point gets control from the cold start loader and is responsible for basic system initialization, including sending a signon message (which can be omitted in the first version). If the IOBYTE function is implemented, it must be set at this point. The various system parameters which are set by the WBOOT entry point must be initialized, and control is transferred to the CCP at 3400H+b for further processing. Note that reg C must be set to zero to select drive A.

WBOOT The WBOOT entry point gets control when a warm start occurs. A warm start is performed whenever a user program branches to location 0000H, or when the CPU is reset from the front panel. The CP/M system must be loaded from the first two tracks of drive A up to, but not including, the BIOS (or CBIOS, if you have completed your patch). System parameters must be initialized as shown below:

location 0,1,2 set to JMP WBOOT for warm starts
(0000H: JMP 4A03H+b)
location 3 set initial value of IOBYTE, if
implemented in your CBIOS
location 5,6,7 set to JMP BDOS, which is the
primary entry point to CP/M for
transient programs. (0005H: JMP
3C06H+b)

(see Section 9 for complete details of page zero use)
Upon completion of the initialization, the WBOOT program must branch to the CCP at 3400H+b to (re)start the system. Upon entry to the CCP, register C is set to the drive to select after system initialization.

CONST Sample the status of the currently assigned console device and return 0FFH in register A if a character is ready to read, and 00H in register A if no console characters are ready.

CONIN Read the next console character into register A, and

set the parity bit (high order bit) to zero. If no console character is ready, wait until a character is typed before returning.

- CONOUT Send the character from register C to the console output device. The character is in ASCII, with high order parity bit set to zero. You may want to include a time-out on a line feed or carriage return, if your console device requires some time interval at the end of the line (such as a TI Silent 700 terminal). You can, if you wish, filter out control characters which cause your console device to react in a strange way (a control-z causes the Lear Seigler terminal to clear the screen, for example).
- LIST Send the character from register C to the currently assigned listing device. The character is in ASCII with zero parity.
- PUNCH Send the character from register C to the currently assigned punch device. The character is in ASCII with zero parity.
- READER Read the next character from the currently assigned reader device into register A with zero parity (high order bit must be zero), an end of file condition is reported by returning an ASCII control-z (1AH).
- HOME Return the disk head of the currently selected disk (initially disk A) to the track 00 position. If your controller allows access to the track 0 flag from the drive, step the head until the track 0 flag is detected. If your controller does not support this feature, you can translate the HOME call into a call on SETTRK with a parameter of 0.
- SELDSK Select the disk drive given by register C for further operations, where register C contains 0 for drive A, 1 for drive B, and so-forth up to 15 for drive P (the standard CP/M distribution version supports four drives). On each disk select, SELDSK must return in HL the base address of a 16-byte area, called the Disk Parameter Header, described in the Section 10. For standard floppy disk drives, the contents of the header and associated tables does not change, and thus the program segment included in the sample CBIOS performs this operation automatically. If there is an attempt to select a non-existent drive, SELDSK returns HL=0000H as an error indicator. Although SELDSK must return the header address on each call, it is advisable to postpone the actual physical disk select operation until an I/O function (seek, read or write) is actually performed, since disk selects often occur without ultimately performing any disk I/O, and many controllers will unload the head of the current disk

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before selecting the new drive. This would cause an excessive amount of noise and disk wear.

SETTRK Register BC contains the track number for subsequent disk accesses on the currently selected drive. You can choose to seek the selected track at this time, or delay the seek until the next read or write actually occurs. Register BC can take on values in the range 0-76 corresponding to valid track numbers for standard floppy disk drives, and 0-65535 for non-standard disk subsystems.

SETSEC Register BC contains the sector number (1 through 26) for subsequent disk accesses on the currently selected drive. You can choose to send this information to the controller at this point, or instead delay sector selection until a read or write operation occurs.

SETDMA Register BC contains the DMA (disk memory access) address for subsequent read or write operations. For example, if B = 00H and C = 80H when SETDMA is called, then all subsequent read operations read their data into 80H through 0FFH, and all subsequent write operations get their data from 80H through 0FFH, until the next call to SETDMA occurs. The initial DMA address is assumed to be 80H. Note that the controller need not actually support direct memory access. If, for example, all data is received and sent through I/O ports, the CBIOS which you construct will use the 128 byte area starting at the selected DMA address for the memory buffer during the following read or write operations.

READ Assuming the drive has been selected, the track has been set, the sector has been set, and the DMA address has been specified, the READ subroutine attempts to read one sector based upon these parameters, and returns the following error codes in register A:

0	no errors occurred
1	non-recoverable error condition occurred

Currently, CP/M responds only to a zero or non-zero value as the return code. That is, if the value in register A is 0 then CP/M assumes that the disk operation completed properly. If an error occurs, however, the CBIOS should attempt at least 10 retries to see if the error is recoverable. When an error is reported the BDOS will print the message "BDOS ERR ON x: BAD SECTOR". The operator then has the option of typing <cr> to ignore the error, or ctl-C to abort.

WRITE Write the data from the currently selected DMA address to the currently selected drive, track, and sector. The data should be marked as "non deleted data" to

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maintain compatibility with other CP/M systems. The error codes given in the READ command are returned in register A, with error recovery attempts as described above.

LISTST Return the ready status of the list device. Used by the DESPOOL program to improve console response during its operation. The value 00 is returned in A if the list device is not ready to accept a character, and 0FFH if a character can be sent to the printer. Note that a 00 value always suffices.

SECTRAN Performs sector logical to physical sector translation in order to improve the overall response of CP/M. Standard CP/M systems are shipped with a "skew factor" of 6, where six physical sectors are skipped between each logical read operation. This skew factor allows enough time between sectors for most programs to load their buffers without missing the next sector. In particular computer systems which use fast processors, memory, and disk subsystems, the skew factor may be changed to improve overall response. Note, however, that you should maintain a single density IBM compatible version of CP/M for information transfer into and out of your computer system, using a skew factor of 6. In general, SECTRAN receives a logical sector number in BC, and a translate table address in DE. The sector number is used as an index into the translate table, with the resulting physical sector number in HL. For standard systems, the tables and indexing code is provided in the CBIOS and need not be changed.

7. A SAMPLE BIOS

The program shown in Appendix C can serve as a basis for your first BIOS. The simplest functions are assumed in this BIOS, so that you can enter it through the front panel, if absolutely necessary. Note that the user must alter and insert code into the subroutines for CONST, CONIN, CONOUT, READ, WRITE, and WAITIO subroutines. Storage is reserved for user-supplied code in these regions. The scratch area reserved in page zero (see Section 9) for the BIOS is used in this program, so that it could be implemented in ROM, if desired.

Once operational, this skeletal version can be enhanced to print the initial sign-on message and perform better error recovery. The subroutines for LIST, PUNCH, and READER can be filled-out, and the IOBYTE function can be implemented.

8. A SAMPLE COLD START LOADER

The program shown in Appendix D can serve as a basis for your cold start loader. The disk read function must be supplied by the user, and the program must be loaded somehow starting at location 0000. Note that space is reserved for your patch so that the total amount of storage required for the cold start loader is 128 bytes. Eventually, you will probably want to get this loader onto the first disk sector (track 0, sector 1), and cause your controller to load it into memory automatically upon system start-up. Alternatively, you may wish to place the cold start loader into ROM, and place it above the CP/M system. In this case, it will be necessary to originate the program at a higher address, and key-in a jump instruction at system start-up which branches to the loader. Subsequent warm starts will not require this key-in operation, since the entry point 'WBOOT' gets control, thus bringing the system in from disk automatically. Note also that the skeletal cold start loader has minimal error recovery, which may be enhanced on later versions.

9. RESERVED LOCATIONS IN PAGE ZERO

Main memory page zero, between locations 00H and 0FFH, contains several segments of code and data which are used during CP/M processing. The code and data areas are given below for reference purposes.

Locations from to	Contents
0000H - 0002H	Contains a jump instruction to the warm start entry point at location 4A03H+b. This allows a simple programmed restart (JMP 0000H) or manual restart from the front panel.
0003H - 0003H	Contains the Intel standard IOBYTE, which is optionally included in the user's CBIOS, as described in Section 6.
0004H - 0004H	Current default drive number (0=A,...,15=P).
0005H - 0007H	Contains a jump instruction to the BDOS, and serves two purposes: JMP 0005H provides the primary entry point to the BDOS, as described in the manual "CP/M Interface Guide," and LHL 0006H brings the address field of the instruction to the HL register pair. This value is the lowest address in memory used by CP/M (assuming the CCP is being overlaid). Note that the DDT program will change the address field to reflect the reduced memory size in debug mode.
0008H - 0027H	(interrupt locations 1 through 5 not used)
0030H - 0037H	(interrupt location 6, not currently used - reserved)
0038H - 003AH	Restart 7 - Contains a jump instruction into the DDT or SID program when running in debug mode for programmed breakpoints, but is not otherwise used by CP/M.
003BH - 003FH	(not currently used - reserved)
0040H - 004FH	16 byte area reserved for scratch by CBIOS, but is not used for any purpose in the distribution version of CP/M
0050H - 005BH	(not currently used - reserved)
005CH - 007CH	default file control block produced for a transient program by the Console Command Processor.
007DH - 007FH	Optional default random record position

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0080H - 00FFH default 128 byte disk buffer (also filled with the command line when a transient is loaded under the CCP).

Note that this information is set-up for normal operation under the CP/M system, but can be overwritten by a transient program if the BDOS facilities are not required by the transient.

If, for example, a particular program performs only simple I/O and must begin execution at location 0, it can be first loaded into the TPA, using normal CP/M facilities, with a small memory move program which gets control when loaded (the memory move program must get control from location 0100H, which is the assumed beginning of all transient programs). The move program can then proceed to move the entire memory image down to location 0, and pass control to the starting address of the memory load. Note that if the BIOS is overwritten, or if location 0 (containing the warm start entry point) is overwritten, then the programmer must bring the CP/M system back into memory with a cold start sequence.

10. DISK PARAMETER TABLES.

Tables are included in the BIOS which describe the particular characteristics of the disk subsystem used with CP/M. These tables can be either hand-coded, as shown in the sample CBIOS in Appendix C, or automatically generated using the DISKDEF macro library, as shown in Appendix B. The purpose here is to describe the elements of these tables.

In general, each disk drive has an associated (16-byte) disk parameter header which both contains information about the disk drive and provides a scratchpad area for certain BDOS operations. The format of the disk parameter header for each drive is shown below

Disk Parameter Header																	
	XLT		0000		0000		0000		DIRBUF		DPB		CSV		ALV		
	16b		16b		16b		16b		16b		16b		16b		16b		16b

where each element is a word (16-bit) value. The meaning of each Disk Parameter Header (DPH) element is

XLT	Address of the logical to physical translation vector, if used for this particular drive, or the value 0000H if no sector translation takes place (i.e, the physical and logical sector numbers are the same). Disk drives with identical sector skew factors share the same translate tables.
0000	Scratchpad values for use within the BDOS (initial value is unimportant).
DIRBUF	Address of a 128 byte scratchpad area for directory operations within BDOS. All DPH's address the same scratchpad area.
DPB	Address of a disk parameter block for this drive. Drives with identical disk characteristics address the same disk parameter block.
CSV	Address of a scratchpad area used for software check for changed disks. This address is different for each DPH.
ALV	Address of a scratchpad area used by the BDOS to keep disk storage allocation information. This address is different for each DPH.

Given n disk drives, the DPH's are arranged in a table whose first row of 16 bytes corresponds to drive 0, with the last row corresponding to drive n-1. The table thus appears as

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DPBASE:

```

-----
00 |XLT 00| 0000 | 0000 | 0000 |DIRBUF|DBP 00|CSV 00|ALV 00|
-----
01 |XLT 01| 0000 | 0000 | 0000 |DIRBUF|DBP 01|CSV 01|ALV 01|
-----
                                (and so-forth through)
-----
n-1|XLTn-1| 0000 | 0000 | 0000 |DIRBUF|DBPn-1|CSVn-1|ALVn-1|
-----

```

where the label DPBASE defines the base address of the DPH table.

A responsibility of the SELDSK subroutine is to return the base address of the DPH for the selected drive. The following sequence of operations returns the table address, with a 0000H returned if the selected drive does not exist.

```

NDISKS EQU 4 ;NUMBER OF DISK DRIVES
*****
SELDISK:
;SELECT DISK GIVEN BY BC
LXI H,0000H ;ERROR CODE
MOV A,C ;DRIVE OK?
CPI NDISKS ;CY IF SO
RNC ;RET IF ERROR
;NO ERROR, CONTINUE
MOV L,C ;LOW(DISK)
MOV H,B ;HIGH(DISK)
DAD H ;*2
DAD H ;*4
DAD H ;*8
DAD H ;*16
LXI D,DPBASE ;FIRST DPH
DAD D ;DPH(DISK)
RET

```

The translation vectors (XLT 00 through XLTn-1) are located elsewhere in the BIOS, and simply correspond one-for-one with the logical sector numbers zero through the sector count-1. The Disk Parameter Block (DPB) for each drive is more complex. A particular DPB, which is addressed by one or more DPH's, takes the general form

```

-----
| SPT |BSH|BLM|EXM| DSM | DRM |AL0|AL1| CKS | OFF |
-----
16b 8b 8b 8b 16b 16b 8b 8b 16b 16b
-----

```

where each is a byte or word value, as shown by the "8b" or "16b" indicator below the field.

SPT is the total number of sectors per track

BSH is the data allocation block shift factor, determined by the data block allocation size.

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EXM is the extent mask, determined by the data block allocation size and the number of disk blocks.

DSM determines the total storage capacity of the disk drive

DRM determines the total number of directory entries which can be stored on this drive AL0,AL1 determine reserved directory blocks.

CKS is the size of the directory check vector

OFF is the number of reserved tracks at the beginning of the (logical) disk.

The values of BSH and BLM determine (implicitly) the data allocation size BLS, which is not an entry in the disk parameter block. Given that the designer has selected a value for BLS, the values of BSH and BLM are shown in the table below

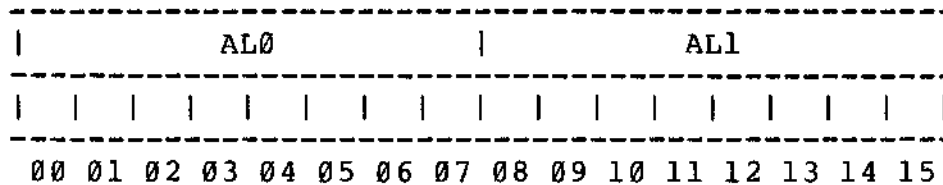
BLS	BSH	BLM
1,024	3	7
2,048	4	15
4,096	5	31
8,192	6	63
16,384	7	127

where all values are in decimal. The value of EXM depends upon both the BLS and whether the DSM value is less than 256 or greater than 255, as shown in the following table

BLS	DSM < 256	DSM > 255
1,024	0	N/A
2,048	1	0
4,096	3	1
8,192	7	3
16,384	15	7

The value of DSM is the maximum data block number supported by this particular drive, measured in BLS units. The product BLS times (DSM+1) is the total number of bytes held by the drive and, of course, must be within the capacity of the physical disk, not counting the reserved operating system tracks.

The DRM entry is the one less than the total number of directory entries, which can take on a 16-bit value. The values of AL0 and AL1, however, are determined by DRM. The two values AL0 and AL1 can together be considered a string of 16-bits, as shown below.



where position 00 corresponds to the high order bit of the byte labelled AL0, and 15 corresponds to the low order bit of the byte labelled AL1. Each bit position reserves a data block for number of directory entries, thus allowing a total of 16 data blocks to be assigned for directory entries (bits are assigned starting at 00 and filled to the right until position 15). Each directory entry occupies 32 bytes, resulting in the following table

BLS	Directory Entries
1,024	32 times # bits
2,048	64 times # bits
4,096	128 times # bits
8,192	256 times # bits
16,384	512 times # bits

Thus, if DRM = 127 (128 directory entries), and BLS = 1024, then there are 32 directory entries per block, requiring 4 reserved blocks. In this case, the 4 high order bits of AL0 are set, resulting in the values AL0 = 0F0H and AL1 = 00H.

The CKS value is determined as follows: if the disk drive media is removable, then $CKS = (DRM+1)/4$, where DRM is the last directory entry number. If the media is fixed, then set $CKS = 0$ (no directory records are checked in this case).

Finally, the OFF field determines the number of tracks which are skipped at the beginning of the physical disk. This value is automatically added whenever SETTRK is called, and can be used as a mechanism for skipping reserved operating system tracks, or for partitioning a large disk into smaller segmented sections.

To complete the discussion of the DPB, recall that several DPH's can address the same DPB if their drive characteristics are identical. Further, the DPB can be dynamically changed when a new drive is addressed by simply changing the pointer in the DPH since the BDOS copies the DPB values to a local area whenever the SELDSK function is invoked.

Returning back to the DPH for a particular drive, note that the two address values CSV and ALV remain. Both addresses reference an area of uninitialized memory following the BIOS. The areas must be unique for each drive, and the size of each area is determined by the values in the DPB.

The size of the area addressed by CSV is CKS bytes, which is sufficient to hold the directory check information for this particular drive. If $CKS = (DRM+1)/4$, then you must reserve $(DRM+1)/4$ bytes for directory check use. If $CKS = 0$, then no storage is reserved.

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The size of the area addressed by ALV is determined by the maximum number of data blocks allowed for this particular disk, and is computed as $(DSM/8)+1$.

The CBIOS shown in Appendix C demonstrates an instance of these tables for standard 8" single density drives. It may be useful to examine this program, and compare the tabular values with the definitions given above.

11. THE DISKDEF MACRO LIBRARY.

A macro library is shown in Appendix F, called DISKDEF, which greatly simplifies the table construction process. You must have access to the MAC macro assembler, of course, to use the DISKDEF facility, while the macro library is included with all CP/M 2.0 distribution disks.

A BIOS disk definition consists of the following sequence of macro statements:

```
MACLIB   DISKDEF
.....
DISKS    n
DISKDEF  0,...
DISKDEF  1,...
.....
DISKDEF  n-1
.....
ENDEF
```

where the MACLIB statement loads the DISKDEF.LIB file (on the same disk as your BIOS) into MAC's internal tables. The DISKS macro call follows, which specifies the number of drives to be configured with your system, where n is an integer in the range 1 to 16. A series of DISKDEF macro calls then follow which define the characteristics of each logical disk, 0 through n-1 (corresponding to logical drives A through P). Note that the DISKS and DISKDEF macros generate the in-line fixed data tables described in the previous section, and thus must be placed in a non-executable portion of your BIOS, typically directly following the BIOS jump vector.

The remaining portion of your BIOS is defined following the DISKDEF macros, with the ENDEF macro call immediately preceding the END statement. The ENDEF (End of Diskdef) macro generates the necessary uninitialized RAM areas which are located in memory above your BIOS.

The form of the DISKDEF macro call is

```
DISKDEF  dn,fsc,lsc,[skf],bls,dks,dir,cks,ofs,[0]
```

where

dn	is the logical disk number, 0 to n-1
fsc	is the first physical sector number (0 or 1)
lsc	is the last sector number
skf	is the optional sector skew factor
bls	is the data allocation block size
dir	is the number of directory entries
cks	is the number of "checked" directory entries
ofs	is the track offset to logical track 00
[0]	is an optional 1.4 compatibility flag

The value "dn" is the drive number being defined with this DISKDEF

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macro invocation. The "fsc" parameter accounts for differing sector numbering systems, and is usually 0 or 1. The "lsc" is the last numbered sector on a track. When present, the "skf" parameter defines the sector skew factor which is used to create a sector translation table according to the skew. If the number of sectors is less than 256, a single-byte table is created, otherwise each translation table element occupies two bytes. No translation table is created if the skf parameter is omitted (or equal to 0). The "bls" parameter specifies the number of bytes allocated to each data block, and takes on the values 1024, 2048, 4096, 8192, or 16384. Generally, performance increases with larger data block sizes since there are fewer directory references and logically connected data records are physically close on the disk. Further, each directory entry addresses more data and the BIOS-resident ram space is reduced. The "dks" specifies the total disk size in "bls" units. That is, if the bls = 2048 and dks = 1000, then the total disk capacity is 2,048,000 bytes. If dks is greater than 255, then the block size parameter bls must be greater than 1024. The value of "dir" is the total number of directory entries which may exceed 255, if desired. The "cks" parameter determines the number of directory items to check on each directory scan, and is used internally to detect changed disks during system operation, where an intervening cold or warm start has not occurred (when this situation is detected, CP/M automatically marks the disk read/only so that data is not subsequently destroyed). As stated in the previous section, the value of cks = dir when the media is easily changed, as is the case with a floppy disk subsystem. If the disk is permanently mounted, then the value of cks is typically 0, since the probability of changing disks without a restart is quite low. The "ofs" value determines the number of tracks to skip when this particular drive is addressed, which can be used to reserve additional operating system space or to simulate several logical drives on a single large capacity physical drive. Finally, the [0] parameter is included when file compatibility is required with versions of 1.4 which have been modified for higher density disks. This parameter ensures that only 16K is allocated for each directory record, as was the case for previous versions. Normally, this parameter is not included.

For convenience and economy of table space, the special form

```
DISKDEF i,j
```

gives disk i the same characteristics as a previously defined drive j. A standard four-drive single density system, which is compatible with version 1.4, is defined using the following macro invocations:

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```

DISKS      4
DISKDEF   0,1,26,6,1024,243,64,64,2
DISKDEF   1,0
DISKDEF   2,0
DISKDEF   3,0
.....
ENDEF

```

with all disks having the same parameter values of 26 sectors per track (numbered 1 through 26), with 6 sectors skipped between each access, 1024 bytes per data block, 243 data blocks for a total of 243k byte disk capacity, 64 checked directory entries, and two operating system tracks.

The DISKS macro generates n Disk Parameter Headers (DPH's), starting at the DPH table address DPBASE generated by the macro. Each disk header block contains sixteen bytes, as described above, and correspond one-for-one to each of the defined drives. In the four drive standard system, for example, the DISKS macro generates a table of the form:

```

DPBASE EQU $
DPE0:  DW  XLT0,0000H,0000H,0000H,DIRBUF,DPB0,CSV0,ALV0
DPE1:  DW  XLT0,0000H,0000H,0000H,DIRBUF,DPB0,CSV1,ALV1
DPE2:  DW  XLT0,0000H,0000H,0000H,DIRBUF,DPB0,CSV2,ALV2
DPE3:  DW  XLT0,0000H,0000H,0000H,DIRBUF,DPB0,CSV3,ALV3

```

where the DPH labels are included for reference purposes to show the beginning table addresses for each drive 0 through 3. The values contained within the disk parameter header are described in detail in the previous section. The check and allocation vector addresses are generated by the ENDEF macro in the ram area following the BIOS code and tables.

Note that if the "skf" (skew factor) parameter is omitted (or equal to 0), the translation table is omitted, and a 0000H value is inserted in the XLT position of the disk parameter header for the disk. In a subsequent call to perform the logical to physical translation, SECTRAN receives a translation table address of DE = 0000H, and simply returns the original logical sector from BC in the HL register pair. A translate table is constructed when the skf parameter is present, and the (non-zero) table address is placed into the corresponding DPH's. The table shown below, for example, is constructed when the standard skew factor skf = 6 is specified in the DISKDEF macro call:

```

XLT0:  DB  1,7,13,19,25,5,11,17,23,3,9,15,21
        DB  2,8,14,20,26,6,12,18,24,4,10,16,22

```

Following the ENDEF macro call, a number of uninitialized data areas are defined. These data areas need not be a part of the BIOS which is loaded upon cold start, but must be available between the BIOS and the end of memory. The size of the uninitialized RAM area is determined by EQU statements generated by the ENDEF macro. For a standard four-drive system, the ENDEF macro might produce

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```

4C72 =      BEGDAT EQU $
          (data areas)
4DB0 =      ENDDAT EQU $
013C =      DATSIZ EQU $-BEGDAT

```

which indicates that uninitialized RAM begins at location 4C72H, ends at 4DB0H-1, and occupies 013CH bytes. You must ensure that these addresses are free for use after the system is loaded.

After modification, you can use the STAT program to check your drive characteristics, since STAT uses the disk parameter block to decode the drive information. The STAT command form

```
STAT d:DSK:
```

decodes the disk parameter block for drive d (d=A,...,P) and displays the values shown below:

```

r: 128 Byte Record Capacity
k: Kilobyte Drive Capacity
d: 32 Byte Directory Entries
c: Checked Directory Entries
e: Records/ Extent
b: Records/ Block
s: Sectors/ Track
t: Reserved Tracks

```

Three examples of DISKDEF macro invocations are shown below with corresponding STAT parameter values (the last produces a full 8-megabyte system).

```

DISKDEF 0,1,58,,2048,256,128,128,2
r=4096, k=512, d=128, c=128, e=256, b=16, s=58, t=2

```

```

DISKDEF 0,1,58,,2048,1024,300,0,2
r=16384, k=2048, d=300, c=0, e=128, b=16, s=58, t=2

```

```

DISKDEF 0,1,58,,16384,512,128,128,2
r=65536, k=8192, d=128, c=128, e=1024, b=128, s=58, t=2

```

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12. SECTOR BLOCKING AND DEBLOCKING.

Upon each call to the BIOS WRITE entry point, the CP/M BDOS includes information which allows effective sector blocking and deblocking where the host disk subsystem has a sector size which is a multiple of the basic 128-byte unit. The purpose here is to present a general-purpose algorithm which can be included within your BIOS which uses the BDOS information to perform the operations automatically.

Upon each call to WRITE, the BDOS provides the following information in register C:

0	=	normal sector write
1	=	write to directory sector
2	=	write to the first sector of a new data block

Condition 0 occurs whenever the next write operation is into a previously written area, such as a random mode record update, when the write is to other than the first sector of an unallocated block, or when the write is not into the directory area. Condition 1 occurs when a write into the directory area is performed. Condition 2 occurs when the first record (only) of a newly allocated data block is written. In most cases, application programs read or write multiple 128 byte sectors in sequence, and thus there is little overhead involved in either operation when blocking and deblocking records since pre-read operations can be avoided when writing records.

Appendix G lists the blocking and deblocking algorithms in skeletal form (this file is included on your CP/M disk). Generally, the algorithms map all CP/M sector read operations onto the host disk through an intermediate buffer which is the size of the host disk sector. Throughout the program, values and variables which relate to the CP/M sector involved in a seek operation are prefixed by "sek," while those related to the host disk system are prefixed by "hst." The equate statements beginning on line 29 of Appendix G define the mapping between CP/M and the host system, and must be changed if other than the sample host system is involved.

The entry points BOOT and WBOOT must contain the initialization code starting on line 57, while the SELDSK entry point must be augmented by the code starting on line 65. Note that although the SELDSK entry point computes and returns the Disk Parameter Header address, it does not physically selected the host disk at this point (it is selected later at READHST or WRITEHST). Further, SETTRK, SETTRK, and SETDMA simply store the values, but do not take any other action at this point. SECTRAN performs a trivial trivial function of returning the physical sector number.

The principal entry points are READ and WRITE, starting on lines 110 and 125, respectively. These subroutines take the place of your previous READ and WRITE operations.

The actual physical read or write takes place at either WRITEHST or READHST, where all values have been prepared: hstdsk is the host

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disk number, hsttrk is the host track number, and hstsec is the host sector number (which may require translation to a physical sector number). You must insert code at this point which performs the full host sector read or write into, or out of, the buffer at hstbuf of length hstsiz. All other mapping functions are performed by the algorithms.

This particular algorithm was tested using an 80 megabyte hard disk unit which was originally configured for 128 byte sectors, producing approximately 35 megabytes of formatted storage. When configured for 512 byte host sectors, usable storage increased to 57 megabytes, with a corresponding 400% improvement in overall response. In this situation, there is no apparent overhead involved in deblocking sectors, with the advantage that user programs still maintain the (less memory consuming) 128-byte sectors. This is primarily due, of course, to the information provided by the BDOS which eliminates the necessity for pre-read operations to take place.

APPENDIX A: THE MDS COLD START LOADER

```

;      MDS-800 Cold Start Loader for CP/M 2.0
;
;      Version 2.0 August, 1979
;
0000 = false equ 0
ffff = true  equ not false
0000 = testing equ false
;
      if testing
bias equ 03400h
      endif
0000 = bias equ 0000h
      endif
0000 = cpmb equ bias ;base of dos load
0806 = bdos equ 806h+bias ;entry to dos for calls
1880 = bdose equ 1880h+bias ;end of dos load
1600 = boot equ 1600h+bias ;cold start entry point
1603 = rboot equ boot+3 ;warm start entry point
;
3000 ; org 3000h ;loaded here by hardware
;
1880 = bdosl equ bdose-cpmb
0002 = ntrks equ 2 ;tracks to read
0031 = bdoss equ bdosl/128 ;# sectors in bdos
0019 = bdos0 equ 25 ;# on track 0
0018 = bdosl equ bdoss-bdos0 ;# on track 1
;
f800 = mon80 equ 0f800h ;intel monitor base
ff0f = rmon80 equ 0ff0fh ;restart location for mon80
0078 = base equ 078h ;'base' used by controller
0079 = rtype equ base+1 ;result type
007b = rbyte equ base+3 ;result byte
007f = reset equ base+7 ;reset controller
;
0078 = dstat equ base ;disk status port
0079 = ilow equ base+1 ;low iopb address
007a = ihigh equ base+2 ;high iopb address
00ff = bsw equ 0ffh ;boot switch
0003 = recal equ 3h ;recalibrate selected drive
0004 = readf equ 4h ;disk read function
0100 = stack equ 100h ;use end of boot for stack
;
rstart:
3000 310001 lxi sp,stack;in case of call to mon80
; clear disk status
3003 db79 in rtype
3005 db7b in rbyte
; check if boot switch is off
coldstart:
3007 dbff in bsw
3009 e602 ani 02h ;switch on?
300b c20730 jnz coldstart

```



```

; clear the controller
300e d37f out reset ;logic cleared
;
;
3010 0602 mvi b,ntrks ;number of tracks to read
3012 214230 lxi h,iopb0
;
start:
;
; read first/next track into cpmb
3015 7d mov a,1
3016 d379 out ilow
3018 7c mov a,h
3019 d37a out ihigh
301b db78 wait0: in dstat
301d e604 ani 4
301f calb30 jz wait0
;
; check disk status
3022 db79 in rtype
3024 e603 ani 11b
3026 fe02 cpi 2
;
; if testing
; cnc rmon80 ;go to monitor if 11 or 10
; endif
; if not testing
3028 d20030 jnc rstart ;retry the load
; endif
;
; in rbyte ;i/o complete, check status
; if not ready, then go to mon80
302b db7b in rbyte
302d 17 ral
302e dc0fff cc rmon80 ;not ready bit set
3031 1f rar ;restore
3032 e61e ani 11110b ;overrun/addr err/seek/crc
;
; if testing
; cnz rmon80 ;go to monitor
; endif
; if not testing
3034 c20030 jnz rstart ;retry the load
; endif
;
;
; lxi d,iopbl ;length of iopb
3037 110700 lxi d,iopbl
303a 19 dad d ;addressing next iopb
303b 05 dcr b ;count down tracks
303c c21530 jnz start
;
;
; jmp boot, print message, set-up jmps
303f c30016 jmp boot
;
; parameter blocks

```

```

3042 80      iopb0: db      80h      ;iocw, no update
3043 04              db      readf    ;read function
3044 19              db      bdos0    ;# sectors to read trk 0
3045 00              db      0        ;track 0
3046 02              db      2        ;start with sector 2, trk 0
3047 0000         dw      cpmb      ;start at base of bdos
0007 =          iopb1 equ      $-iopb0
;
3049 80      iopb1: db      80h
304a 04              db      readf
304b 18              db      bdos1    ;sectors to read on track 1
304c 01              db      1        ;track 1
304d 01              db      1        ;sector 1
304e 800c         dw      cpmb+bdos0*128 ;base of second rd
3050          end

```

APPENDIX B: THE MDS BASIC I/O SYSTEM (BIOS)

```

;      mds-800 i/o drivers for cp/m 2.0
;      (four drive single density version)
;
;      version 2.0 august, 1979
;
0014 = vers      equ      20      ;version 2.0
;
;      copyright (c) 1979
;      digital research
;      box 579, pacific grove
;      california, 93950
;
4a00      org      4a00h      ;base of bios in 20k system
3400 =    cpmb     equ      3400h      ;base of cpm ccp
3c06 =    bdos     equ      3c06h      ;base of bdos in 20k system
1600 =    cpml     equ      $-cpmb     ;length (in bytes) of cpm system
002c =    nsects  equ      cpml/128;number of sectors to load
0002 =    offset  equ      2          ;number of disk tracks used by cp
0004 =    cdisk   equ      0004h      ;address of last logged disk
0080 =    buff    equ      0080h      ;default buffer address
000a =    retry   equ      10         ;max retries on disk i/o before e
;
;      perform following functions
;      boot      cold start
;      wboot     warm start (save i/o byte)
;      (boot and wboot are the same for mds)
;      const     console status
;              reg-a = 00 if no character ready
;              reg-a = ff if character ready
;      conin     console character in (result in reg-a)
;      conout    console character out (char in reg-c)
;      list      list out (char in reg-c)
;      punch     punch out (char in reg-c)
;      reader    paper tape reader in (result to reg-a)
;      home      move to track 00
;
;      (the following calls set-up the io parameter bloc
;      mds, which is used to perform subsequent reads an
;      seldsk   select disk given by reg-c (0,1,2,...)
;      settrk   set track address (0,...76) for sub r/w
;      setsec   set sector address (1,...,26)
;      setdma   set subsequent dma address (initially 80h)
;
;      read/write assume previous calls to set i/o parms
;      read     read track/sector to preset dma address
;      write    write track/sector from preset dma address
;
;      jump vector for individual routines
4a00 c3b34a      jmp      boot
4a03 c3c34a wboote: jmp      wboot
4a06 c3614b      jmp      const
4a09 c3644b      jmp      conin
4a0c c36a4b      jmp      conout

```

```

4a0f c36d4b      jmp      list
4a12 c3724b      jmp      punch
4a15 c3754b      jmp      reader
4a18 c3784b      jmp      home
4a1b c37d4b      jmp      seldsk
4a1e c3a74b      jmp      settrk
4a21 c3ac4b      jmp      setsec
4a24 c3bb4b      jmp      setdma
4a27 c3c14b      jmp      read
4a2a c3ca4b      jmp      write
4a2d c3704b      jmp      listst ;list status
4a30 c3b14b      jmp      sectran
;
;          maclib  diskdef ;load the disk definition library
disks      4          ;four disks
4a33+=     dpbase   equ      $          ;base of disk parameter blocks
4a33+824a00 dpe0:    dw      xlt0,0000h      ;translate table
4a37+000000 dw      0000h,0000h      ;scratch area
4a3b+6e4c73 dw      dirbuf,dpb0        ;dir buff,param block
4a3f+0d4dee dw      csv0,alv0          ;check, alloc vectors
4a43+824a00 dpe1:    dw      xlt1,0000h      ;translate table
4a47+000000 dw      0000h,0000h      ;scratch area
4a4b+6e4c73 dw      dirbuf,dpb1        ;dir buff,param block
4a4f+3c4d1d dw      csv1,alv1          ;check, alloc vectors
4a53+824a00 dpe2:    dw      xlt2,0000h      ;translate table
4a57+000000 dw      0000h,0000h      ;scratch area
4a5b+6e4c73 dw      dirbuf,dpb2        ;dir buff,param block
4a5f+6b4d4c dw      csv2,alv2          ;check, alloc vectors
4a63+824a00 dpe3:    dw      xlt3,0000h      ;translate table
4a67+000000 dw      0000h,0000h      ;scratch area
4a6b+6e4c73 dw      dirbuf,dpb3        ;dir buff,param block
4a6f+9a4d7b dw      csv3,alv3          ;check, alloc vectors
diskdef    0,1,26,6,1024,243,64,64,offset
4a73+=     dpb0     equ      $          ;disk parm block
4a73+1a00  dw      26          ;sec per track
4a75+03    db      3          ;block shift
4a76+07    db      7          ;block mask
4a77+00    db      0          ;extnt mask
4a78+f200  dw      242        ;disk size-1
4a7a+3f00  dw      63         ;directory max
4a7c+c0    db      192       ;alloc0
4a7d+00    db      0         ;allocl
4a7e+1000  dw      16        ;check size
4a80+0200  dw      2         ;offset
4a82+=     xlt0     equ      $          ;translate table
4a82+01    db      1         ;
4a83+07    db      7         ;
4a84+0d    db      13        ;
4a85+13    db      19        ;
4a86+19    db      25        ;
4a87+05    db      5         ;
4a88+0b    db      11        ;
4a89+11    db      17        ;
4a8a+17    db      23        ;
4a8b+03    db      3         ;

```

```

4a8c+09      db      9
4a8d+0f      db      15
4a8e+15      db      21
4a8f+02      db      2
4a90+08      db      8
4a91+0e      db      14
4a92+14      db      20
4a93+1a      db      26
4a94+06      db      6
4a95+0c      db      12
4a96+12      db      18
4a97+18      db      24
4a98+04      db      4
4a99+0a      db      10
4a9a+10      db      16
4a9b+16      db      22
              diskdef 1,0
4a73+=       dpb1    equ    dpb0    ;equivalent parameters
001f+=       als1    equ    als0    ;same allocation vector size
0010+=       css1    equ    css0    ;same checksum vector size
4a82+=       xlt1    equ    xlt0    ;same translate table
              diskdef 2,0
4a73+=       dpb2    equ    dpb0    ;equivalent parameters
001f+=       als2    equ    als0    ;same allocation vector size
0010+=       css2    equ    css0    ;same checksum vector size
4a82+=       xlt2    equ    xlt0    ;same translate table
              diskdef 3,0
4a73+=       dpb3    equ    dpb0    ;equivalent parameters
001f+=       als3    equ    als0    ;same allocation vector size
0010+=       css3    equ    css0    ;same checksum vector size
4a82+=       xlt3    equ    xlt0    ;same translate table
              ;
              ;      undef occurs at end of assembly
              ;
              ;      end of controller - independent code, the remaini
              ;      are tailored to the particular operating environm
              ;      be altered for any system which differs from the
              ;
              ;      the following code assumes the mds monitor exists
              ;      and uses the i/o subroutines within the monitor
              ;
              ;      we also assume the mds system has four disk drive
00fd =       revrt   equ    0fdh    ;interrupt revert port
00fc =       intc    equ    0fch    ;interrupt mask port
00f3 =       icon    equ    0f3h    ;interrupt control port
007e =       inte    equ    0111$1110b;enable rst 0(warm boot),rst 7
              ;
              ;      mds monitor equates
f800 =       mon80   equ    0f800h  ;mds monitor
ff0f =       rmon80 equ    0ff0fh   ;restart mon80 (boot error)
f803 =       ci      equ    0f803h  ;console character to reg-a
f806 =       ri      equ    0f806h  ;reader in to reg-a
f809 =       co      equ    0f809h  ;console char from c to console o
f80c =       po      equ    0f80ch  ;punch char from c to punch devic
f80f =       lo      equ    0f80fh  ;list from c to list device
f812 =       csts    equ    0f812h  ;console status 00/ff to register

```

```

;
; disk ports and commands
0078 = base equ 78h ;base of disk command io ports
0078 = dstat equ base ;disk status (input)
0079 = rtype equ base+1 ;result type (input)
007b = rbyte equ base+3 ;result byte (input)
;
0079 = ilow equ base+1 ;iopb low address (output)
007a = ihigh equ base+2 ;iopb high address (output)
;
0004 = readf equ 4h ;read function
0006 = writf equ 6h ;write function
0003 = recal equ 3h ;recalibrate drive
0004 = iordy equ 4h ;i/o finished mask
000d = cr equ 0dh ;carriage return
000a = lf equ 0ah ;line feed
;
signon: ;signon message: xxk cp/m vers y.y
4a9c 0d0a0a db cr,lf,lf
4a9f 3230 db '20' ;sample memory size
4aa1 6b2043f db 'k cp/m vers '
4aad 322e30 db vers/10+'0','.',vers mod 10+'0'
4ab0 0d0a00 db cr,lf,0
;
boot: ;print signon message and go to ccp
; (note: mds boot initialized iobyte at 0003h)
4ab3 310001 lxi sp,buff+80h
4ab6 219c4a lxi h,signon
4ab9 cdd34b call prmsg ;print message
4abc af xra a ;clear accumulator
4abd 320400 sta cdisk ;set initially to disk a
4ac0 c30f4b jmp gocpm ;go to cp/m
;
;
wboot:; loader on track 0, sector 1, which will be skippe
; read cp/m from disk - assuming there is a 128 byt
; start.
;
4ac3 318000 lxi sp,buff ;using dma - thus 80 thru ff ok f
;
4ac6 0e0a mvi c,retry ;max retries
4ac8 c5 push b
wboot0: ;enter here on error retries
4ac9 010034 lxi b,cpmb ;set dma address to start of disk
4acc cdbb4b call setdma
4acf 0e00 mvi c,0 ;boot from drive 0
4ad1 cd7d4b call seldisk
4ad4 0e00 mvi c,0
4ad6 cda74b call settrk ;start with track 0
4ad9 0e02 mvi c,2 ;start reading sector 2
4adb cdac4b call setsec
;
; read sectors, count nsects to zero
4ade cl pop b ;10-error count
4adf 062c mvi b,nsects

```

```

rdsec: ;read next sector
4ae1 c5      push    b          ;save sector count
4ae2 cdcl4b  call    read
4ae5 c2494b  jnz    booterr ;retry if errors occur
4ae8 2a6c4c  lhld   iod          ;increment dma address
4aeb 118000  lxi    d,128       ;sector size
4aee 19      dad    d          ;incremented dma address in hl
4aef 44      mov    b,h
4af0 4d      mov    c,l          ;ready for call to set dma
4af1 cdbb4b  call   setdma
4af4 3a6b4c  lda    ios          ;sector number just read
4af7 fela    cpi    26          ;read last sector?
4af9 da054b  jc     rd1
;          must be sector 26, zero and go to next track
4afc 3a6a4c  lda    iot          ;get track to register a
4aff 3c      inr    a
4b00 4f      mov    c,a          ;ready for call
4b01 cda74b  call   settrk
4b04 af      xra    a          ;clear sector number
4b05 3c      rdl:   inr    a          ;to next sector
4b06 4f      mov    c,a          ;ready for call
4b07 cdac4b  call   setsec
4b0a cl     pop    b          ;recall sector count
4b0b 05     dcr    b          ;done?
4b0c c2e14a  jnz   rdsec
;
;          done with the load, reset default buffer address
gocpm: ;(enter here from cold start boot)
;          enable rst0 and rst7
4b0f f3     di
4b10 3e12   mvi    a,12h       ;initialize command
4b12 d3fd   out    revrt
4b14 af     xra    a
4b15 d3fc   out    intc        ;cleared
4b17 3e7e   mvi    a,inte      ;rst0 and rst7 bits on
4b19 d3fc   out    intc
4b1b af     xra    a
4b1c d3f3   out    icon        ;interrupt control
;
;          set default buffer address to 80h
4b1e 018000 lxi    b,buff
4b21 cdbb4b  call   setdma
;
;          reset monitor entry points
4b24 3ec3   mvi    a,jmp
4b26 320000 sta    0
4b29 21034a lxi    h,wboote
4b2c 220100 shld   1           ;jmp wboot at location 00
4b2f 320500 sta    5
4b32 21063c lxi    h,bdos
4b35 220600 shld   6           ;jmp bdos at location 5
4b38 323800 sta    7*8        ;jmp to mon80 (may have been chan
4b3b 2100f8 lxi    h,mon80
4b3e 223900 shld   7*8+1
;          leave iobyte set

```

```

;          previously selected disk was b, send parameter to
4b41 3a0400      lda      cdisk      ;last logged disk number
4b44 4f          mov      c,a        ;send to ccp to log it in
4b45 fb          ei
4b46 c30034      jmp      cpmb
;
;          error condition occurred, print message and retry
booterr:
4b49 c1          pop      b          ;recall counts
4b4a 0d          dcr      c
4b4b ca524b      jz       booter0
;          try again
4b4e c5          push     b
4b4f c3c94a      jmp      wboot0
;
booter0:
;          otherwise too many retries
4b52 215b4b      lxi     h,bootmsg
4b55 cdd34b      call    prmsg
4b58 c30fff      jmp     rmon80 ;mds hardware monitor
;
bootmsg:
4b5b 3f626f4     db      '?boot',0
;
;
const: ;console status to reg-a
;      (exactly the same as mds call)
4b61 c312f8     jmp     csts
;
conin: ;console character to reg-a
4b64 cd03f8     call    ci
4b67 e67f       ani     7fh      ;remove parity bit
4b69 c9         ret
;
conout: ;console character from c to console out
4b6a c309f8     jmp     co
;
list: ;list device out
;      (exactly the same as mds call)
4b6d c30ff8     jmp     lo
;
listst:
;return list status
4b70 af         xra     a
4b71 c9         ret          ;always not ready
;
punch: ;punch device out
;      (exactly the same as mds call)
4b72 c30cf8     jmp     po
;
reader: ;reader character in to reg-a
;      (exactly the same as mds call)
4b75 c306f8     jmp     ri
;
home: ;move to home position

```



```

;      treat as track 00 seek
4b78 0e00      mvi      c,0
4b7a c3a74b    jmp      settrk
;
; seldsk: ;select disk given by register c
4b7d 210000    lxi      h,0000h ;return 0000 if error
4b80 79        mov      a,c
4b81 fe04      cpi      ndisks ;too large?
4b83 d0        rnc      ;leave hl = 0000
;
4b84 e602      ani      l0b     ;00 00 for drive 0,1 and 10 10 fo
4b86 32664c    sta      dbank   ;to select drive bank
4b89 79        mov      a,c     ;00, 01, 10, 11
4b8a e601      ani      lb      ;mds has 0,1 at 78, 2,3 at 88
4b8c b7        ora      a       ;result 00?
4b8d ca924b    jz      setdrive
4b90 3e30      mvi      a,00110000b ;selects drive 1 in bank
setdrive:
4b92 47        mov      b,a     ;save the function
4b93 21684c    lxi      h,iof   ;io function
4b96 7e        mov      a,m
4b97 e6cf      ani      l1001111b ;mask out disk number
4b99 b0        ora      b       ;mask in new disk number
4b9a 77        mov      m,a     ;save it in iopb
4b9b 69        mov      l,c
4b9c 2600      mvi      h,0     ;hl=disk number
4b9e 29        dad      h       ;*2
4b9f 29        dad      h       ;*4
4ba0 29        dad      h       ;*8
4ba1 29        dad      h       ;*16
4ba2 11334a    lxi      d,dibase
4ba5 19        dad      d       ;hl=disk header table address
4ba6 c9        ret
;
;
; settrk: ;set track address given by c
4ba7 216a4c    lxi      h,iot
4baa 71        mov      m,c
4bab c9        ret
;
; setsec: ;set sector number given by c
4bac 216b4c    lxi      h,ios
4baf 71        mov      m,c
4bb0 c9        ret
sectran:
;translate sector bc using table at de
4bb1 0600      mvi      b,0     ;double precision sector number i
4bb3 eb      xchg      ;translate table address to hl
4bb4 09        dad      b       ;translate(sector) address
4bb5 7e        mov      a,m     ;translated sector number to a
4bb6 326b4c    sta      ios
4bb9 6f        mov      l,a
4bba c9        ret
;
; setdma: ;set dma address given by regs b,c

```

```

4bbb 69          mov     l,c
4bbc 60          mov     h,b
4bbd 226c4c     shld   iod
4bc0 c9          ret

;
read:           ;read next disk record (assuming disk/trk/sec/dma
4bc1 0e04       mvi     c,readf ;set to read function
4bc3 cde04b     call   setfunc
4bc6 cdf04b     call   waitio  ;perform read function
4bc9 c9          ret                ;may have error set in reg-a

;
;
write:         ;disk write function
4bca 0e06       mvi     c,writf
4bcc cde04b     call   setfunc ;set to write function
4bcf cdf04b     call   waitio
4bd2 c9          ret                ;may have error set

;
;
; utility subroutines
prmsg:        ;print message at h,l to 0
4bd3 7e          mov     a,m
4bd4 b7          ora     a          ;zero?
4bd5 c8          rz

; more to print
4bd6 e5          push   h
4bd7 4f          mov     c,a
4bd8 cd6a4b     call   conout
4bdb e1          pop    h
4bdc 23          inx    h
4bdd c3d34b     jmp    prmsg

;
setfunc:      ;
; set function for next i/o (command in reg-c)
4be0 21684c     lxi    h,iof    ;io function address
4be3 7e          mov     a,m      ;get it to accumulator for maskin
4be4 e6f8       ani    11111000b ;remove previous command
4be6 b1          ora     c        ;set to new command
4be7 77          mov     m,a      ;replaced in iopb
; the mds-800 controller reg's disk bank bit in sec
; mask the bit from the current i/o function
4be8 e620       ani    00100000b ;mask the disk select bit
4bea 216b4c     lxi    h,ios    ;address the sector selec
4bed b6          ora     m        ;select proper disk bank
4bee 77          mov     m,a      ;set disk select bit on/o
4bef c9          ret

;
waitio:       ;
4bf0 0e0a       mvi     c,retry ;max retries before perm error
rewait:      ;
; start the i/o function and wait for completion
4bf2 cd3f4c     call   intype  ;in rtype
4bf5 cd4c4c     call   inbyte  ;clears the controller

;
4bf8 3a664c     lda     dbank   ;set bank flags

```

```

4bfb b7          ora      a          ;zero if drive 0,1 and nz
4bfc 3e67       mvi      a,iopb and 0ffh ;low address for iopb
4bfe 064c       mvi      b,iopb shr 8    ;high address for iopb
4c00 c20b4c     jnz      iodrl          ;drive bank 1?
4c03 d379       out      ilow          ;low address to controlle
4c05 78         mov      a,b
4c06 d37a       out      ihigh         ;high address
4c08 c3104c     jmp      wait0         ;to wait for complete
;
iodrl:          ;drive bank 1
4c0b d389       out      ilow+10h      ;88 for drive bank 10
4c0d 78         mov      a,b
4c0e d38a       out      ihigh+10h
;
4c10 cd594c     wait0:   call     instat        ;wait for completion
4c13 e604       ani      iordy         ;ready?
4c15 ca104c     jz       wait0
;
;          check io completion ok
4c18 cd3f4c     call     intype        ;must be io complete (00)
;          00 unlinked i/o complete, 01 linked i/o comple
;          10 disk status changed 11 (not used)
4c1b fe02       cpi      l0b          ;ready status change?
4c1d ca324c     jz       wready
;
;          must be 00 in the accumulator
4c20 b7         ora      a
4c21 c2384c     jnz      werror        ;some other condition, re
;
;          check i/o error bits
4c24 cd4c4c     call     inbyte
4c27 17         ral
4c28 da324c     jc       wready        ;unit not ready
4c2b 1f         rar
4c2c e6fe       ani      l1111110b    ;any other errors?
4c2e c2384c     jnz      werror
;
;          read or write is ok, accumulator contains zero
4c31 c9         ret
;
wready:        ;not ready, treat as error for now
4c32 cd4c4c     call     inbyte        ;clear result byte
4c35 c3384c     jmp      trycount
;
werror:        ;return hardware malfunction (crc, track, seek, e
;          the mds controller has returned a bit in each pos
;          of the accumulator, corresponding to the conditio
;          0 - deleted data (accepted as ok above)
;          1 - crc error
;          2 - seek error
;          3 - address error (hardware malfunction)
;          4 - data over/under flow (hardware malfunct
;          5 - write protect (treated as not ready)
;          6 - write error (hardware malfunction)
;          7 - not ready

```

```

;      (accumulator bits are numbered 7 6 5 4 3 2 1 0)
;
;      it may be useful to filter out the various condit
;      but we will get a permanent error message if it i
;      recoverable.  in any case, the not ready conditio
;      treated as a separate condition for later improve
trycount:
;      register c contains retry count, decrement 'til z
4c38 0d      dcr      c
4c39 c2f24b  jnz      await ;for another try
;
;      cannot recover from error
4c3c 3e01    mvi      a,1      ;error code
4c3e c9      ret
;
;      intype, inbyte, instat read drive bank 00 or 10
4c3f 3a664c intype: lda      dbank
4c42 b7      ora      a
4c43 c2494c jnz      intypl ;skip to bank 10
4c46 db79    in       rtype
4c48 c9      ret
4c49 db89    intypl: in      rtype+10h      ;78 for 0,1 88 for 2,3
4c4b c9      ret
;
4c4c 3a664c inbyte: lda      dbank
4c4f b7      ora      a
4c50 c2564c jnz      inbytl
4c53 db7b    in       rbyte
4c55 c9      ret
4c56 db8b    inbytl: in      rbyte+10h
4c58 c9      ret
;
4c59 3a664c instat: lda      dbank
4c5c b7      ora      a
4c5d c2634c jnz      instal
4c60 db78    in       dstat
4c62 c9      ret
4c63 db88    instal: in      dstat+10h
4c65 c9      ret
;
;
;      data areas (must be in ram)
4c66 00      dbank: db      0      ;disk bank 00 if drive 0,1
;                  ;                  10 if drive 2,3
;
;      iopb: ;io parameter block
4c67 80      db      80h     ;normal i/o operation
4c68 04      ioof: db      readf  ;io function, initial read
4c69 01      ion:  db      1      ;number of sectors to read
4c6a 02      iot:  db      offset ;track number
4c6b 01      ios:  db      1      ;sector number
4c6c 8000    iod:  dw      buff   ;io address
;
;
;      define ram areas for bdos operation

```

```

                                .def      .endef
4c6e+=      begdat  equ      $
4c6e+      dirbuf: ds      128      ;directory access buffer
4ceet+      alv0:   ds      31
4d0d+      csv0:   ds      16
4d1d+      alv1:   ds      31
4d3c+      csv1:   ds      16
4d4c+      alv2:   ds      31
4d6b+      csv2:   ds      16
4d7b+      alv3:   ds      31
4d9a+      csv3:   ds      16
4daa+=      enddat  equ      $
013c+=      datsiz  equ      $-begdat
4daa      end

```

APPENDIX C: A SKELETAL CBIOS

```

; skeletal cbios for first level of cp/m 2.0 altera
;
0014 = msize equ 20 ;cp/m version memory size in kilo
;
; "bias" is address offset from 3400h for memory sy
; than 16k (referred to as "b" throughout the text)
;
0000 = bias equ (msize-20)*1024
3400 = ccp equ 3400h+bias ;base of ccp
3c06 = bdos equ ccp+806h ;base of bdos
4a00 = bios equ ccp+1600h ;base of bios
0004 = cdisk equ 0004h ;current disk number 0=a,...,15=p
0003 = iobyte equ 0003h ;intel i/o byte
;
4a00 org bios ;origin of this program
002c = nsects equ ($-ccp)/128 ;warm start sector count
;
; jump vector for individual subroutines
4a00 c39c4a jmp boot ;cold start
4a03 c3a64a wboote: jmp wboot ;warm start
4a06 c3114b jmp const ;console status
4a09 c3244b jmp conin ;console character in
4a0c c3374b jmp conout ;console character out
4a0f c3494b jmp list ;list character out
4a12 c3464b jmp punch ;punch character out
4a15 c34f4b jmp reader ;reader character out
4a18 c3544b jmp home ;move head to home positi
4a1b c35a4b jmp seldisk ;select disk
4a1e c37d4b jmp settrk ;set track number
4a21 c3924b jmp setsec ;set sector number
4a24 c3ad4b jmp setdma ;set dma address
4a27 c3c34b jmp read ;read disk
4a2a c3d64b jmp write ;write disk
4a2d c34b4b jmp listst ;return list status
4a30 c3a74b jmp sectran ;sector translate
;
; fixed data tables for four-drive standard
; ibm-compatible 8" disks
; disk parameter header for disk 00
4a33 734a00 dpbase: dw trans,0000h
4a37 000000 dw 0000h,0000h
4a3b f04c8d dw dirbf,dpblk
4a3f ec4d70 dw chk00,all00
; disk parameter header for disk 01
4a43 734a00 dw trans,0000h
4a47 000000 dw 0000h,0000h
4a4b f04c8d dw dirbf,dpblk
4a4f fc4d8f dw chk01,all01
; disk parameter header for disk 02
4a53 734a00 dw trans,0000h
4a57 000000 dw 0000h,0000h
4a5b f04c8d dw dirbf,dpblk
4a5f 0c4eae dw chk02,all02

```

```

;          disk parameter header for disk 03
4a63 734a00      dw      trans,0000h
4a67 000000      dw      0000h,0000h
4a6b f04c8d      dw      dirbf,dpblk
4a6f 1c4ecd      dw      chk03,a1103

;
;          sector translate vector
trans:
4a73 01070d      db      1,7,13,19      ;sectors 1,2,3,4
4a77 19050b      db      25,5,11,17      ;sectors 5,6,7,8
4a7b 170309      db      23,3,9,15      ;sectors 9,10,11,12
4a7f 150208      db      21,2,8,14      ;sectors 13,14,15,16
4a83 141a06      db      20,26,6,12      ;sectors 17,18,19,20
4a87 121804      db      18,24,4,10      ;sectors 21,22,23,24
4a8b 1016        db      16,22          ;sectors 25,26

;
dpblk: ;disk parameter block, common to all disks
4a8d 1a00        dw      26          ;sectors per track
4a8f 03          db      3          ;block shift factor
4a90 07          db      7          ;block mask
4a91 00          db      0          ;null mask
4a92 f200        dw      242         ;disk size-1
4a94 3f00        dw      63          ;directory max
4a96 c0          db      192         ;alloc 0
4a97 00          db      0          ;alloc 1
4a98 1000        dw      16          ;check size
4a9a 0200        dw      2          ;track offset

;
;          end of fixed tables
;
;          individual subroutines to perform each function
boot: ;simplest case is to just perform parameter initi
4a9c af          xra      a          ;zero in the accum
4a9d 320300      sta      iobyte      ;clear the iobyte
4aa0 320400      sta      cdisk       ;select disk zero
4aa3 c3ef4a      jmp      gocpm        ;initialize and go to cp/

;
wboot: ;simplest case is to read the disk until all sect
4aa6 318000      lxi      sp,80h      ;use space below buffer f
4aa9 0e00        mvi      c,0         ;select disk 0
4aab cd5a4b      call     seldsk      ;select disk
4aae cd544b      call     home        ;go to track 00

;
4ab1 062c        mvi      b,nsects   ;b counts # of sectors to
4ab3 0e00        mvi      c,0         ;c has the current track
4ab5 1602        mvi      d,2         ;d has the next sector to

;          note that we begin by reading track 0, sector 2 s
;          contains the cold start loader, which is skipped
4ab7 210034      lxi      h,ccp       ;base of cp/m (initial lo

loadl: ;load one more sector
4aba c5          push     b          ;save sector count, current track
4abb d5          push     d          ;save next sector to read
4abc e5          push     h          ;save dma address
4abd 4a          mov      c,d         ;get sector address to register c
4abe cd924b      call     setsec     ;set sector address from register
4acl c1          pop      b          ;recall dma address to b,c

```

```

4ac2 c5      push    b      ;replace on stack for later recal
4ac3 cdad4b  call    setdma ;set dma address from b,c
;
; drive set to 0, track set, sector set, dma address
4ac6 cdc34b  call    read
4ac9 fe00    cpi     00h    ;any errors?
4acb c2a64a  jnz    wboot  ;retry the entire boot if an erro
;
; no error, move to next sector
4ace e1      pop     h      ;recall dma address
4acf 118000  lxi    d,128  ;dma=dma+128
4ad2 19     dad    d      ;new dma address is in h,l
4ad3 d1      pop     d      ;recall sector address
4ad4 c1      pop     b      ;recall number of sectors remaini
4ad5 05     dcr    b      ;sectors=sectors-1
4ad6 caef4a  jz     gocpm  ;transfer to cp/m if all have bee
;
; more sectors remain to load, check for track chan
4ad9 14     inr    d
4ada 7a     mov    a,d    ;sector=27?, if so, change tracks
4adb felb   cpi    27
4add daba4a  jc     loadl  ;carry generated if sector<27
;
; end of current track, go to next track
4ae0 1601   mvi    d,1    ;begin with first sector of next
4ae2 0c     inr    c      ;track=track+1
;
; save register state, and change tracks
4ae3 c5     push   b
4ae4 d5     push   d
4ae5 e5     push   h
4ae6 cd7d4b call   settrk ;track address set from register
4ae9 e1     pop    h
4aea d1     pop    d
4aeb c1     pop    b
4aec c3ba4a jmp    loadl  ;for another sector
;
; end of load operation, set parameters and go to c
gocpm:
4aef 3ec3   mvi    a,0c3h ;c3 is a jmp instruction
4af1 320000 sta    0      ;for jmp to wboot
4af4 21034a lxi    h,wboote ;wboot entry point
4af7 220100 shld   l      ;set address field for jmp at 0
;
4afa 320500 sta    5      ;for jmp to bdos
4afd 21063c lxi    h,bdos ;bdos entry point
4b00 220600 shld   6      ;address field of jump at 5 to bd
;
4b03 018000 lxi    b,80h  ;default dma address is 80h
4b06 cdad4b call   setdma
;
4b09 fb     ei        ;enable the interrupt system
4b0a 3a0400 lda    cdisk  ;get current disk number
4b0d 4f     mov    c,a   ;send to the ccp
4b0e c30034 jmp    ccp    ;go to cp/m for further processin

```



```

;
;
;   simple i/o handlers (must be filled in by user)
;   in each case, the entry point is provided, with s
;   to insert your own code
;
const: ;console status, return 0ffh if character ready,
4b11      ds      10h      ;space for status subroutine
4b21 3e00  mvi      a,00h
4b23 c9    ret

;
conin:  ;console character into register a
4b24      ds      10h      ;space for input routine
4b34 e67f  ani      7fh      ;strip parity bit
4b36 c9    ret

;
conout: ;console character output from register c
4b37 79    mov      a,c      ;get to accumulator
4b38      ds      10h      ;space for output routine
4b48 c9    ret

;
list:   ;list character from register c
4b49 79    mov      a,c      ;character to register a
4b4a c9    ret          ;null subroutine

;
listst: ;return list status (0 if not ready, 1 if ready)
4b4b af    xra      a      ;0 is always ok to return
4b4c c9    ret

;
punch:  ;punch character from register c
4b4d 79    mov      a,c      ;character to register a
4b4e c9    ret          ;null subroutine

;
;
reader: ;read character into register a from reader device
4b4f 3ela  mvi      a,lah      ;enter end of file for now (repla
4b51 e67f  ani      7fh      ;remember to strip parity bit
4b53 c9    ret

;
;
;   i/o drivers for the disk follow
;   for now, we will simply store the parameters away
;   in the read and write subroutines
;
home:   ;move to the track 00 position of current drive
;       translate this call into a settrk call with param
4b54 0e00  mvi      c,0      ;select track 0
4b56 cd7d4b call     settrk
4b59 c9    ret          ;we will move to 00 on first read

;
seldsk: ;select disk given by register c
4b5a 210000 lxi     h,0000h ;error return code
4b5d 79    mov      a,c
4b5e 32ef4c sta     diskno
4b61 fe04  cpi      4      ;must be between 0 and 3

```

```

4b63 d0          rnc          ;no carry if 4,5,...
;
4b64             ds          10h ;space for disk select
;
4b6e 3aef4c     lda          diskno
4b71 6f         mov          l,a   ;l=disk number 0,1,2,3
4b72 2600      mvi          h,0   ;high order zero
4b74 29        dad          h     ;*2
4b75 29        dad          h     ;*4
4b76 29        dad          h     ;*8
4b77 29        dad          h     ;*16 (size of each header)
4b78 11334a    lxi          d,dpbase
4b7b 19        dad          d     ;hl=.dpbase(diskno*16)
4b7c c9        ret
;
;
settrk: ;set track given by register c
4b7d 79        mov          a,c
4b7e 32e94c    sta          track
4b81          ds          10h    ;space for track select
4b91 c9        ret
;
;
setsec: ;set sector given by register c
4b92 79        mov          a,c
4b93 32eb4c    sta          sector
4b96          ds          10h    ;space for sector select
4ba6 c9        ret
;
;
sectran:
;translate the sector given by bc using the
;translate table given by de
4ba7 eb        xchg          ;hl=.trans
4ba8 09        dad          b     ;hl=.trans(sector)
4ba9 6e        mov          l,m     ;l = trans(sector)
4baa 2600      mvi          h,0   ;hl= trans(sector)
4bac c9        ret          ;with value in hl
;
;
setdma: ;set dma address given by registers b and c
4bad 69        mov          l,c     ;low order address
4bae 60        mov          h,b     ;high order address
4baf 22ed4c    shld         dmaad   ;save the address
4bb2          ds          10h    ;space for setting the dma address
4bc2 c9        ret
;
;
read: ;perform read operation (usually this is similar
; so we will allow space to set up read command, th
; common code in write)
4bc3          ds          10h    ;set up read command
4bd3 c3e64b    jmp          waitio ;to perform the actual i/o
;
;
write: ;perform a write operation
4bd6          ds          10h    ;set up write command
;
;
waitio: ;enter here from read and write to perform the ac
; operation. return a 00h in register a if the ope
; properly, and 01h if an error occurs during the r

```

```

;
;   in this case, we have saved the disk number in 'd
;   the track number in 'track' (0-76
;   the sector number in 'sector' (1-
;   the dma address in 'dmaad' (0-655
4be6      ds      256      ;space reserved for i/o drivers
4ce6 3e01 mvi      a,1      ;error condition
4ce8 c9    ret          ;replaced when filled-in
;
;   the remainder of the cbios is reserved uninitiali
;   data area, and does not need to be a part of the
;   system memory image (the space must be available,
;   however, between "begdat" and "enddat").
;
4ce9      track: ds      2      ;two bytes for expansion
4ceb      sector: ds     2      ;two bytes for expansion
4ced      dmaad: ds     2      ;direct memory address
4cef      diskno: ds     1      ;disk number 0-15
;
;   scratch ram area for bdos use
4cf0 =    begdat equ     $      ;beginning of data area
4cf0      dirbf: ds    128      ;scratch directory area
4d70      all00: ds     31      ;allocation vector 0
4d8f      all01: ds     31      ;allocation vector 1
4dae      all02: ds     31      ;allocation vector 2
4dcd      all03: ds     31      ;allocation vector 3
4dec      chk00: ds     16      ;check vector 0
4dfc      chk01: ds     16      ;check vector 1
4e0c      chk02: ds     16      ;check vector 2
4e1c      chk03: ds     16      ;check vector 3
;
4e2c =    enddat equ     $      ;end of data area
013c =    datsiz equ    $-begdat;size of data area
4e2c      end

```

APPENDIX D: A SKELETAL GETSYS/PUTSYS PROGRAM

```

;      combined getsys and putsys programs from Sec 4.
;      Start the programs at the base of the TPA

0100          org      0100h

0014 =      msize   equ      20          ; size of cp/m in Kbytes

; "bias" is the amount to add to addresses for > 20k
;      (referred to as "b" throughout the text)

0000 =      bias    equ      (msize-20)*1024
3400 =      ccp     equ      3400h+bias
3c00 =      bdos   equ      ccp+0800h
4a00 =      bios   equ      ccp+1600h

;      getsys programs tracks 0 and 1 to memory at
;      3880h + bias

;      register          usage
;      a                (scratch register)
;      b                track count (0...76)
;      c                sector count (1...26)
;      d,e              (scratch register pair)
;      h,l              load address
;      sp               set to stack address

gstart:
0100 318033    lxi      sp,ccp-0080h      ; start of getsys
0103 218033    lxi      h,ccp-0080h      ; convenient plac
0106 0600      mvi      b,0          ; set initial loa
; start with trac
rd$trk:
0108 0e01      mvi      c,1          ; read next track
; each track star
rd$sec:
010a cd0003    call     read$sec      ; get the next se
010d 118000    lxi      d,128          ; offset by one s
0110 19        dad      d          ; (hl=hl+128)
0111 0c        inr      c          ; next sector
0112 79        mov      a,c          ; fetch sector nu
0113 fe0b      cpi      27          ; and see if la
0115 da0a01    jc       rdsec        ; <, do one more

; arrive here at end of track, move to next track

0118 04        inr      b          ; track = track+1
0119 78        mov      a,b          ; check for last
011a fe02      cpi      2          ; track = 2 ?
011c da0801    jc       rd$trk       ; <, do another

; arrive here at end of load, halt for lack of anything b

011f fb        ei
0120 76        hlt

```

```

;      putsys program, places memory image starting at
;      3880h + bias back to tracks 0 and 1
;      start this program at the next page boundary

0200          org      ($+0100h) and 0ff00h

      put$sys:
0200 318033    lxi      sp,ccp-0080h      ; convenient plac
0203 218033    lxi      h,ccp-0080h      ; start of dump
0206 0600      mvi      b,0              ; start with trac
      wr$trk:
0208 0e01      mvi      c,1              ; start with sect
      wr$sec:
020a cd0004    call     write$sec        ; write one secto
020d 118000    lxi      d,128            ; length of each
0210 19        dad      d                ; <hl>=<hl> + 128
0211 0c        inr     c                  ; <c> = <c> + 1
0212 79        mov     a,c                ; see if
0213 fe0b      cpi     27                 ; past end of t
0215 da0a02    jc      wr$sec            ; no, do another

; arrive here at end of track, move to next track

0218 04        inr     b                  ; track = track+1
0219 78        mov     a,b                ; see if
021a fe02      cpi     2                  ; last track
021c da0802    jc      wr$trk            ; no, do another

; done with putsys, halt for lack of anything better

021f fb        ei
0220 76        hlt

; user supplied subroutines for sector read and write

; move to next page boundary

0300          org      ($+0100h) and 0ff00h

      read$sec:
; read the next sector
; track in <b>,
; sector in <c>
; dmaaddr in <hl>

0300 c5        push    b
0301 e5        push    h

; user defined read operation goes here
0302          ds      64

0342 e1        pop     h
0343 c1        pop     b

```

```

0344 c9          ret
0400          org      ($+0100h) and 0ff00h      ; another page bo
write$sec:
          ; same parameters as read$sec
0400 c5          push   b
0401 e5          push   h
          ; user defined write operation goes here
0402          ds      64
0442 e1          pop    h
0443 c1          pop    b
0444 c9          ret
          ; end of getsys/putsys program
0445          end

```

APPENDIX E: A SKELETAL COLD START LOADER

```

; this is a sample cold start loader which, when modified
; resides on track 00, sector 01 (the first sector on the
; diskette). we assume that the controller has loaded
; this sector into memory upon system start-up (this pro-
; gram can be keyed-in, or can exist in read/only memory
; beyond the address space of the cp/m version you are
; running). the cold start loader brings the cp/m system
; into memory at "loadp" (3400h + "bias"). in a 20k
; memory system, the value of "bias" is 0000h, with large
; values for increased memory sizes (see section 2). afte
; loading the cp/m system, the cold start loader branches
; to the "boot" entry point of the bios, which begins at
; "bios" + "bias." the cold start loader is not used un-
; til the system is powered up again, as long as the bios
; is not overwritten. the origin is assumed at 0000h, an
; must be changed if the controller brings the cold start
; loader into another area, or if a read/only memory area
; is used.

```

```

0000          org      0          ; base of ram in cp/m

0014 =        msize   equ      20          ; min mem size in kbytes

0000 =        bias    equ      (msize-20)*1024 ; offset from 20k system
3400 =        ccp     equ      3400h+bias   ; base of the ccp
4a00 =        bios    equ      ccp+1600h   ; base of the bios
0300 =        biosl   equ      0300h      ; length of the bios
4a00 =        boot    equ      bios
1900 =        size    equ      bios+biosl-ccp ; size of cp/m system
0032 =        sects   equ      size/128    ; # of sectors to load

;          begin the load operation

cold:
0000 010200   lxi      b,2          ; b=0, c=sector 2
0003 1632     mvi      d,sects      ; d=# sectors to load
0005 210034   lxi      h,ccp       ; base transfer address

lsect: ; load the next sector

;          insert inline code at this point to
;          read one 128 byte sector from the
;          track given in register b, sector
;          given in register c,
;          into the address given by <hl>
;
;          branch to location "cold" if a read error occurs

```

```

; *****
; *
; *      user supplied read operation goes here...
; *
; *****

0008 c36b00      jmp      past$patch      ; remove this when patche
000b             ds        60h

      past$patch:
; go to next sector if load is incomplete
006b 15         dcr      d          ; sects=sects-1
006c ca004a     jz       boot           ; head for the bios

;      more sectors to load
;
; we aren't using a stack, so use <sp> as scratch registe
;      to hold the load address increment

006f 318000     lxi     sp,128          ; 128 bytes per sector
0072 39         dad     sp          ; <hl> = <hl> + 128

0073 0c         inr     c          ; sector = sector + 1
0074 79         mov     a,c
0075 felb      cpi     27         ; last sector of track?
0077 da0800     jc      lsect         ; no, go read another

; end of track, increment to next track

007a 0e01     mvi     c,1           ; sector = 1
007c 04         inr     b          ; track = track + 1
007d c30800     jmp     lsect         ; for another group
0080             end       ; of boot loader

```


APPENDIX F: CP/M DISK DEFINITION LIBRARY

```

1: ;      CP/M 2.0 disk re-definition library
2: ;
3: ;      Copyright (c) 1979
4: ;      Digital Research
5: ;      Box 579
6: ;      Pacific Grove, CA
7: ;      93950
8: ;
9: ;      CP/M logical disk drives are defined using the
10: ;     macros given below, where the sequence of calls
11: ;     is:
12: ;
13: ;     disks      n
14: ;     diskdef parameter-list-0
15: ;     diskdef parameter-list-1
16: ;     ...
17: ;     diskdef parameter-list-n
18: ;     endef
19: ;
20: ;     where n is the number of logical disk drives attached
21: ;     to the CP/M system, and parameter-list-i defines the
22: ;     characteristics of the ith drive (i=0,1,...,n-1)
23: ;
24: ;     each parameter-list-i takes the form
25: ;           dn,fsc,lsc,[skf],bls,dks,dir,cks,ofs,[0]
26: ;     where
27: ;     dn      is the disk number 0,1,...,n-1
28: ;     fsc     is the first sector number (usually 0 or 1)
29: ;     lsc     is the last sector number on a track
30: ;     skf     is optional "skew factor" for sector translate
31: ;     bls     is the data block size (1024,2048,...,16384)
32: ;     dks     is the disk size in bls increments (word)
33: ;     dir     is the number of directory elements (word)
34: ;     cks     is the number of dir elements to checksum
35: ;     ofs     is the number of tracks to skip (word)
36: ;     [0]     is an optional 0 which forces 16K/directory en
37: ;
38: ;     for convenience, the form
39: ;           dn,dm
40: ;     defines disk dn as having the same characteristics as
41: ;     a previously defined disk dm.
42: ;
43: ;     a standard four drive CP/M system is defined by
44: ;           disks      4
45: ;           diskdef 0,1,26,6,1024,243,64,64,2
46: ;     dsk      set      0
47: ;           rept      3
48: ;     dsk      set      dsk+1
49: ;           diskdef %dsk,0
50: ;           endm
51: ;           endef
52: ;
53: ;     the value of "begdat" at the end of assembly defines t

```

```

54: ; beginning of the uninitialized ram area above the bios,
55: ; while the value of "enddat" defines the next location
56: ; following the end of the data area. the size of this
57: ; area is given by the value of "datsiz" at the end of t
58: ; assembly. note that the allocation vector will be qui
59: ; large if a large disk size is defined with a small blo
60: ; size.
61: ;
62: dskhdr macro dn
63: ;; define a single disk header list
64: dpe&dn: dw xlt&dn,0000h ;translate table
65: dw 0000h,0000h ;scratch area
66: dw dirbuf,dpb&dn ;dir buff,param block
67: dw csv&dn,alv&dn ;check, alloc vectors
68: endm
69: ;
70: disks macro nd
71: ;; define nd disks
72: ndisks set nd ;;for later reference
73: dpbase equ $ ;base of disk parameter blocks
74: ;; generate the nd elements
75: dsknxt set 0
76: rept nd
77: dskhdr %dsknxt
78: dsknxt set dsknxc+1
79: endm
80: endm
81: ;
82: dpbhdr macro dn
83: dpb&dn equ $ ;disk parm block
84: endm
85: ;
86: ddb macro data,comment
87: ;; define a db statement
88: db data comment
89: endm
90: ;
91: ddw macro data,comment
92: ;; define a dw statement
93: dw data comment
94: endm
95: ;
96: gcd macro m,n
97: ;; greatest common divisor of m,n
98: ;; produces value gcdn as result
99: ;; (used in sector translate table generation)
100: gcdm set m ;;variable for m
101: gcdn set n ;;variable for n
102: gcdr set 0 ;;variable for r
103: rept 65535
104: gcdx set gcdm/gcdn
105: gcdr set gcdm - gcdx*gcdn
106: if gcdr = 0
107: exitm
108: endif

```

```

109: gcdm      set      gcdn
110: gcdn      set      gcdr
111:           endm
112:           endm
113: ;
114: diskdef macro  dn,fsc,lsc,skf,bls,dks,dir,cks,bfs,kl6
115: ;;          generate the set statements for later tables
116:           if      nul lsc
117: ;;          current disk dn same as previous fsc
118: dpb&dn     equ      dpb&fsc ;equivalent parameters
119: als&dn     equ      als&fsc ;same allocation vector size
120: css&dn     equ      css&fsc ;same checksum vector size
121: xlt&dn     equ      xlt&fsc ;same translate table
122:           else
123: secmax     set      lsc-(fsc)      ;;sectors 0...secmax
124: sectors    set      secmax+1;;number of sectors
125: als&dn     set      (dks)/8      ;;size of allocation vector
126:           if      ((dks) mod 8) ne 0
127: als&dn     set      als&dn+1
128:           endif
129: css&dn     set      (cks)/4      ;;number of checksum elements
130: ;;          generate the block shift value
131: blkval     set      bls/128      ;;number of sectors/block
132: blkshf     set      0            ;;counts right 0's in blkval
133: blkmsk     set      0            ;;fills with 1's from right
134:           rept    16            ;;once for each bit position
135:           if      blkval=1
136:           exitm
137:           endif
138: ;;          otherwise, high order 1 not found yet
139: blkshf     set      blkshf+1
140: blkmsk     set      (blkmsk shl 1) or 1
141: blkval     set      blkval/2
142:           endm
143: ;;          generate the extent mask byte
144: blkval     set      bls/1024      ;;number of kilobytes/block
145: extmsk     set      0            ;;fill from right with 1's
146:           rept    16
147:           if      blkval=1
148:           exitm
149:           endif
150: ;;          otherwise more to shift
151: extmsk     set      (extmsk shl 1) or 1
152: blkval     set      blkval/2
153:           endm
154: ;;          may be double byte allocation
155:           if      (dks) > 256
156: extmsk     set      (extmsk shr 1)
157:           endif
158: ;;          may be optional [0] in last position
159:           if      not nul kl6
160: extmsk     set      kl6
161:           endif
162: ;;          now generate directory reservation bit vector
163: dirrem     set      dir          ;;# remaining to process

```

```

164: dirbks  set    bls/32  ;;number of entries per block
165: dirblk  set    0      ;;fill with 1's on each loop
166:         rept    16
167:         if     dirrem=0
168:         exitm
169:         endif
170: ;;      not complete, iterate once again
171: ;;      shift right and add 1 high order bit
172: dirblk  set    (dirblk shr 1) or 8000h
173:         if     dirrem > dirbks
174: dirrem  set    dirrem-dirbks
175:         else
176: dirrem  set    0
177:         endif
178:         endm
179: dpbhdr  dn      ;;generate equ $
180: ddw    %sectors,<;sec per track>
181: ddb    %blkshf,<;block shift>
182: ddb    %blkmsk,<;block mask>
183: ddb    %extmsk,<;extnt mask>
184: ddw    %(dks)-1,<;disk size-1>
185: ddw    %(dir)-1,<;directory max>
186: ddb    %dirblk shr 8,<;alloc0>
187: ddb    %dirblk and 0ffh,<;alloc1>
188: ddw    %(cks)/4,<;check size>
189: ddw    %ofs,<;offset>
190: ;;      generate the translate table, if requested
191:         if     nul skf
192: xlt&dn  equ    0          ;no xlate table
193:         else
194:         if     skf = 0
195: xlt&dn  equ    0          ;no xlate table
196:         else
197: ;;      generate the translate table
198: nxtsec  set    0          ;;next sector to fill
199: nextbas set    0          ;;moves by one on overflow
200:         gcd    %sectors,skf
201: ;;      gcdn = gcd(sectors,skew)
202: neltst  set    sectors/gcdn
203: ;;      neltst is number of elements to generate
204: ;;      before we overlap previous elements
205: nelts  set    neltst  ;;counter
206: xlt&dn  equ    $          ;translate table
207:         rept    sectors ;;once for each sector
208:         if     sectors < 256
209:         ddb    %nxtsec+(fsc)
210:         else
211:         ddw    %nxtsec+(fsc)
212:         endif
213: nxtsec  set    nxtsec+(skf)
214:         if     nxtsec >= sectors
215: nxtsec  set    nxtsec-sectors
216:         endif
217: nelts  set    nelts-1
218:         if     nelts = 0

```

```

219: nxtbas set      nxtbas+1
220: nxtsec set      nxtbas
221: nelts  set      neltst
222:        endif
223:        endm
224:        endif ;;end of nul fac test
225:        endif ;;end of nul bls test
226:        endm
227: ;
228: defds  macro lab,space
229: lab:   ds      space
230:        endm
231: ;
232: lds    macro lb,dn,val
233:        defds lb&dn,%val&dn
234:        endm
235: ;
236: endef  macro
237: ;; generate the necessary ram data areas
238: begdat equ      $
239: dirbuf: ds      128      ;directory access buffer
240: dsknxt set      0
241:        rept ndisks ;;once for each disk
242:        lds   alv,%dsknxt,als
243:        lds   csv,%dsknxt,css
244: dsknxt set      dsknxt+1
245:        endm
246: enddat equ      $
247: datsiz equ      $-begdat
248: ;; db 0 at this point forces hex record
249:        endm

```

APPENDIX G: BLOCKING AND DEBLOCKING ALGORITHMS.

```

1: ;*****
2: ;*
3: ;*      Sector Deblocking Algorithms for CP/M 2.0      *
4: ;*
5: ;*****
6: ;
7: ;      utility macro to compute sector mask
8: smask macro hblk
9: ;;      compute log2(hblk), return @x as result
10: ;;      (2 ** @x = hblk on return)
11: @y set hblk
12: @x set 0
13: ;;      count right shifts of @y until = 1
14:      rept 8
15:      if @y = 1
16:      exitm
17:      endif
18: ;;      @y is not 1, shift right one position
19: @y set @y shr 1
20: @x set @x + 1
21:      endm
22:      endm
23: ;
24: ;*****
25: ;*
26: ;*      CP/M to host disk constants
27: ;*
28: ;*****
29: blksiz equ 2048 ;CP/M allocation size
30: hstsiz equ 512 ;host disk sector size
31: hstspt equ 20 ;host disk sectors/trk
32: hstblk equ hstsiz/128 ;CP/M sects/host buff
33: cpmspt equ hstblk * hstspt ;CP/M sectors/track
34: secmsk equ hstblk-1 ;sector mask
35:      smask hstblk ;compute sector mask
36: secshf equ @x ;log2(hstblk)
37: ;
38: ;*****
39: ;*
40: ;*      BDOS constants on entry to write
41: ;*
42: ;*****
43: wrall equ 0 ;write to allocated
44: wrdir equ 1 ;write to directory
45: wrual equ 2 ;write to unallocated
46: ;
47: ;*****
48: ;*
49: ;*      The BDOS entry points given below show the
50: ;*      code which is relevant to deblocking only.
51: ;*
52: ;*****
53: ;

```

```

54: ; DISKDEF macro, or hand coded tables go here
55: dpbase equ $ ;disk param block base
56: ;
57: boot:
58: wboot:
59: ;enter here on system boot to initialize
60: xra a ;0 to accumulator
61: sta hstact ;host buffer inactive
62: sta unacct ;clear unalloc count
63: ret
64: ;
65: seldsk:
66: ;select disk
67: mov a,c ;selected disk number
68: sta sekdisk ;seek disk number
69: mov l,a ;disk number to HL
70: mvi h,0
71: rept 4 ;multiply by 16
72: dad h
73: endm
74: lxi d,dpbase ;base of parm block
75: dad d ;hl=.dpb(curdisk)
76: ret
77: ;
78: settrk:
79: ;set track given by registers BC
80: mov h,b
81: mov l,c
82: shld sektrk ;track to seek
83: ret
84: ;
85: setsec:
86: ;set sector given by register c
87: mov a,c
88: sta seksec ;sector to seek
89: ret
90: ;
91: setdma:
92: ;set dma address given by BC
93: mov h,b
94: mov l,c
95: shld dmaadr
96: ret
97: ;
98: sectran:
99: ;translate sector number BC
100: mov h,b
101: mov l,c
102: ret
103: ;

```

```

104: ;*****
105: ;*
106: ;*   The READ entry point takes the place of
107: ;*   the previous BIOS definition for READ.
108: ;*
109: ;*****
110: read:
111:     ;read the selected CP/M sector
112:     mvi     a,l
113:     sta     readop           ;read operation
114:     sta     rsflag          ;must read data
115:     mvi     a,wrual
116:     sta     wrtype          ;treat as unalloc
117:     jmp     rwooper         ;to perform the read
118: ;
119: ;*****
120: ;*
121: ;*   The WRITE entry point takes the place of
122: ;*   the previous BIOS definition for WRITE.
123: ;*
124: ;*****
125: write:
126:     ;write the selected CP/M sector
127:     xra     a                ;0 to accumulator
128:     sta     readop          ;not a read operation
129:     mov     a,c              ;write type in c
130:     sta     wrtype
131:     cpi     wrual           ;write unallocated?
132:     jnz     chkuna         ;check for unalloc
133: ;
134: ;   write to unallocated, set parameters
135:     mvi     a,blksiz/128    ;next unalloc recs
136:     sta     unacnt
137:     lda     sekdisk         ;disk to seek
138:     sta     unadsk         ;unadsk = sekdisk
139:     lhld   sektrk
140:     shld   unatrck         ;unatrck = sectrk
141:     lda     seksec
142:     sta     unasec         ;unasec = seksec
143: ;
144: chkuna:
145:     ;check for write to unallocated sector
146:     lda     unacnt         ;any unalloc remain?
147:     ora     a
148:     jz     alloc          ;skip if not
149: ;
150: ;   more unallocated records remain
151:     dcr     a                ;unacnt = unacnt-1
152:     sta     unacnt
153:     lda     sekdisk         ;same disk?
154:     lxi     h,unadsk
155:     cmp     m                ;sekdisk = unadsk?
156:     jnz     alloc          ;skip if not
157: ;
158: ;   disks are the same

```



```

159:      lxi      h,unatrkr
160:      call    sektrkcmp      ;sektrk = unatrkr?
161:      jnz     alloc         ;skip if not
162: ;
163: ;      tracks are the same
164:      lda     seksec         ;same sector?
165:      lxi     h,unasec
166:      cmp     m              ;seksec = unasec?
167:      jnz     alloc         ;skip if not
168: ;
169: ;      match, move to next sector for future ref
170:      inr     m              ;unasec = unasec+1
171:      mov     a,m           ;end of track?
172:      cpi     cpmspt        ;count CP/M sectors
173:      jc      noovf         ;skip if no overflow
174: ;
175: ;      overflow to next track
176:      mvi     m,0           ;unasec = 0
177:      lhld   unatrkr
178:      inx     h
179:      shld   unatrkr       ;unatrkr = unatrkr+1
180: ;
181: noovf:
182:      ;match found, mark as unnecessary read
183:      xra     a             ;0 to accumulator
184:      sta     rsflag        ;rsflag = 0
185:      jmp     rwoper        ;to perform the write
186: ;
187: alloc:
188:      ;not an unallocated record, requires pre-read
189:      xra     a             ;0 to accum
190:      sta     unacnt        ;unacnt = 0
191:      inr     a             ;1 to accum
192:      sta     rsflag        ;rsflag = 1
193: ;
194: ;*****
195: ;*
196: ;*      Common code for READ and WRITE follows
197: ;*
198: ;*****
199: rwoper:
200:      ;enter here to perform the read/write
201:      xra     a             ;zero to accum
202:      sta     erflag        ;no errors (yet)
203:      lda     seksec        ;compute host sector
204:      rept   secshf
205:      ora     a             ;carry = 0
206:      rar     ;shift right
207:      endm
208:      sta     sekfst        ;host sector to seek
209: ;
210: ;      active host sector?
211:      lxi     h,hstact      ;host active flag
212:      mov     a,m
213:      mvi     m,1          ;always becomes 1

```

```

214:      ora      a                ;was it already?
215:      jz       filhst          ;fill host if not
216:      ;
217:      ;      host buffer active, same as seek buffer?
218:      lda      sekdisk
219:      lxi      h,hstdisk        ;same disk?
220:      cmp      m                ;sekdisk = hstdisk?
221:      jnz      nomatch
222:      ;
223:      ;      same disk, same track?
224:      lxi      h,hsttrk
225:      call     sektrkcmp        ;sektrk = hsttrk?
226:      jnz      nomatch
227:      ;
228:      ;      same disk, same track, same buffer?
229:      lda      sekhst
230:      lxi      h,hstsec        ;sekhst = hstsec?
231:      cmp      m
232:      jz       match           ;skip if match
233:      ;
234: nomatch:
235:      ;proper disk, but not correct sector
236:      lda      hstwrtr          ;host written?
237:      ora      a
238:      cnz      writehst        ;clear host buff
239:      ;
240: filhst:
241:      ;may have to fill the host buffer
242:      lda      sekdisk
243:      sta      hstdisk
244:      lhld    sektrk
245:      shld    hsttrk
246:      lda      sekhst
247:      sta      hstsec
248:      lda      rsflag          ;need to read?
249:      ora      a
250:      cnz      readhst         ;yes, if 1
251:      xra      a               ;0 to accum
252:      sta      hstwrtr        ;no pending write
253:      ;
254: match:
255:      ;copy data to or from buffer
256:      lda      seksec          ;mask buffer number
257:      ani      secmsk          ;least signif bits
258:      mov      l,a             ;ready to shift
259:      mvi      h,0             ;double count
260:      rept    7                ;shift left 7
261:      dad     h
262:      endm
263:      ;      hl has relative host buffer address
264:      lxi      d,hstbuf
265:      dad     d                ;hl = host address
266:      xchg
267:      lhld    dmaadr          ;get/put CP/M data
268:      mvi      c,128          ;length of move

```

```

269:      lda      readop          ;which way?
270:      ora      a
271:      jnz      rwmove          ;skip if read
272:      ;
273:      ;      write operation, mark and switch direction
274:      mvi      a,1
275:      sta      hstwrt          ;hstwrt = 1
276:      xchg     ;source/dest swap
277:      ;
278: rwmove:
279:      ;C initially 128, DE is source, HL is dest
280:      ldax    d                ;source character
281:      inx     d
282:      mov     m,a              ;to dest
283:      inx     h
284:      dcr    c                ;loop 128 times
285:      jnz    rwmove
286:      ;
287:      ;      data has been moved to/from host buffer
288:      lda     wrtype           ;write type
289:      cpi    wrdir            ;to directory?
290:      lda     erflag          ;in case of errors
291:      rnz    ;no further processing
292:      ;
293:      ;      clear host buffer for directory write
294:      ora     a                ;errors?
295:      rnz    ;skip if so
296:      xra    a                ;0 to accum
297:      sta    hstwrt           ;buffer written
298:      call   writehst
299:      lda    erflag
300:      ret
301:      ;
302:      ;*****
303:      ;*
304:      ;*      Utility subroutine for 16-bit compare
305:      ;*
306:      ;*****
307: sektrkcmp:
308:      ;HL = .unatrk or .hsttrk, compare with sektrk
309:      xchg
310:      lxi    h,sektrk
311:      ldax  d                ;low byte compare
312:      cmp   m                ;same?
313:      rnz   ;return if not
314:      ;      low bytes equal, test high 1s
315:      inx   d
316:      inx   h
317:      ldax  d
318:      cmp   m                ;sets flags
319:      ret
320:      ;

```

```

321: ;*****
322: ;*
323: ;* WRITEHST performs the physical write to *
324: ;* the host disk, READHST reads the physical *
325: ;* disk. *
326: ;*
327: ;*****
328: writehst:
329: ;hstdsk = host disk #, hsttrk = host track #,
330: ;hstsec = host sect #. write "hstsiz" bytes
331: ;from hstbuf and return error flag in erflag.
332: ;return erflag non-zero if error
333: ret
334: ;
335: readhst:
336: ;hstdsk = host disk #, hsttrk = host track #,
337: ;hstsec = host sect #. read "hstsiz" bytes
338: ;into hstbuf and return error flag in erflag.
339: ret
340: ;
341: ;*****
342: ;*
343: ;* Unitialized RAM data areas *
344: ;*
345: ;*****
346: ;
347: sekdisk: ds 1 ;seek disk number
348: sektrk: ds 2 ;seek track number
349: seksec: ds 1 ;seek sector number
350: ;
351: hstdsk: ds 1 ;host disk number
352: hsttrk: ds 2 ;host track number
353: hstsec: ds 1 ;host sector number
354: ;
355: sekhst: ds 1 ;seek shr secshf
356: hstact: ds 1 ;host active flag
357: hstwrt: ds 1 ;host written flag
358: ;
359: unacct: ds 1 ;unalloc rec cnt
360: unadsk: ds 1 ;last unalloc disk
361: unatrk: ds 2 ;last unalloc track
362: unasec: ds 1 ;last unalloc sector
363: ;
364: erflag: ds 1 ;error reporting
365: rsflag: ds 1 ;read sector flag
366: readop: ds 1 ;1 if read operation
367: wrtype: ds 1 ;write operation type
368: dmaadr: ds 2 ;last dma address
369: hstbuf: ds hstsiz ;host buffer
370: ;

```

```
371: ;*****
372: ;*
373: ;*      The ENDEF macro invocation goes here      *
374: ;*
375: ;*****
376:      end
```

