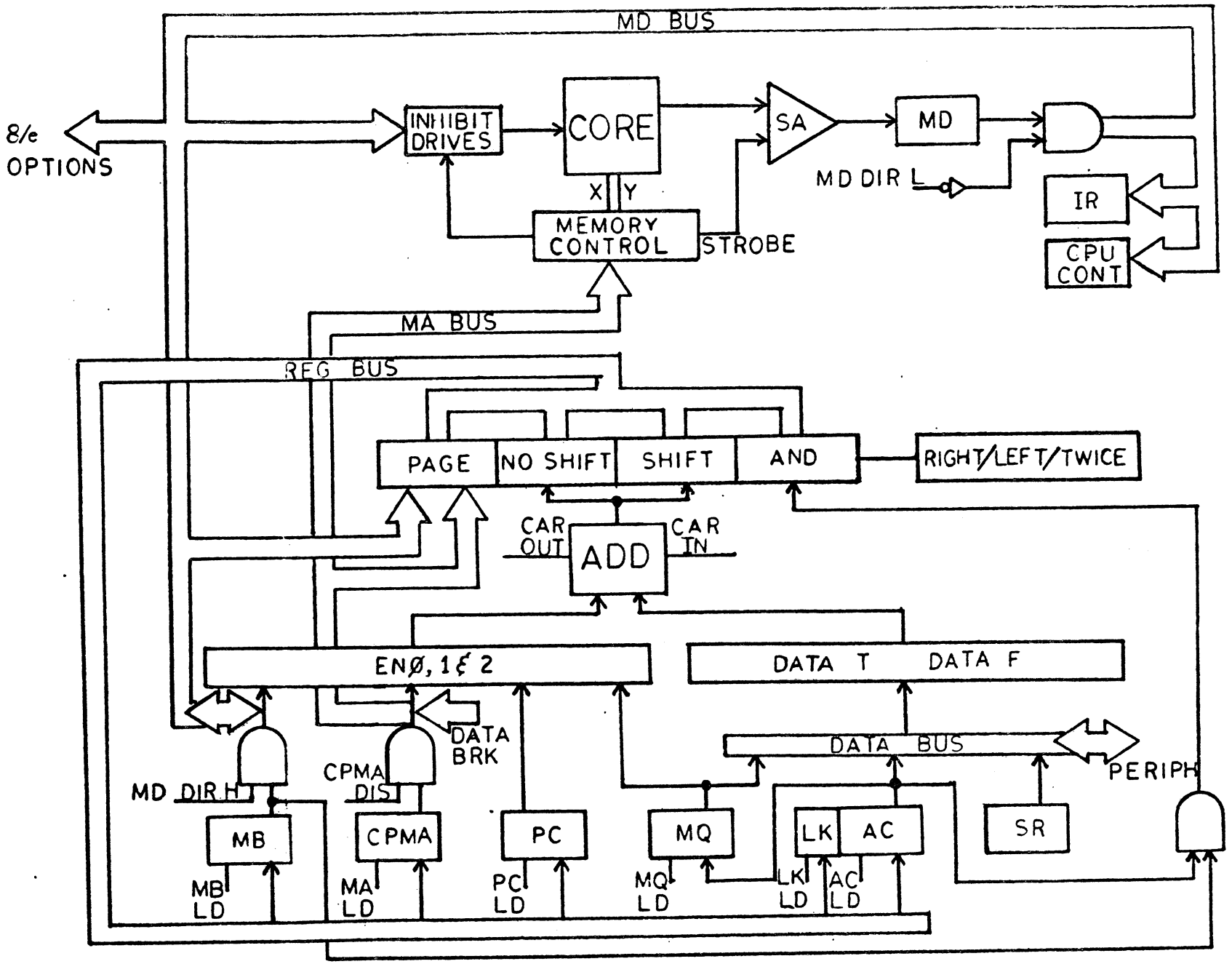
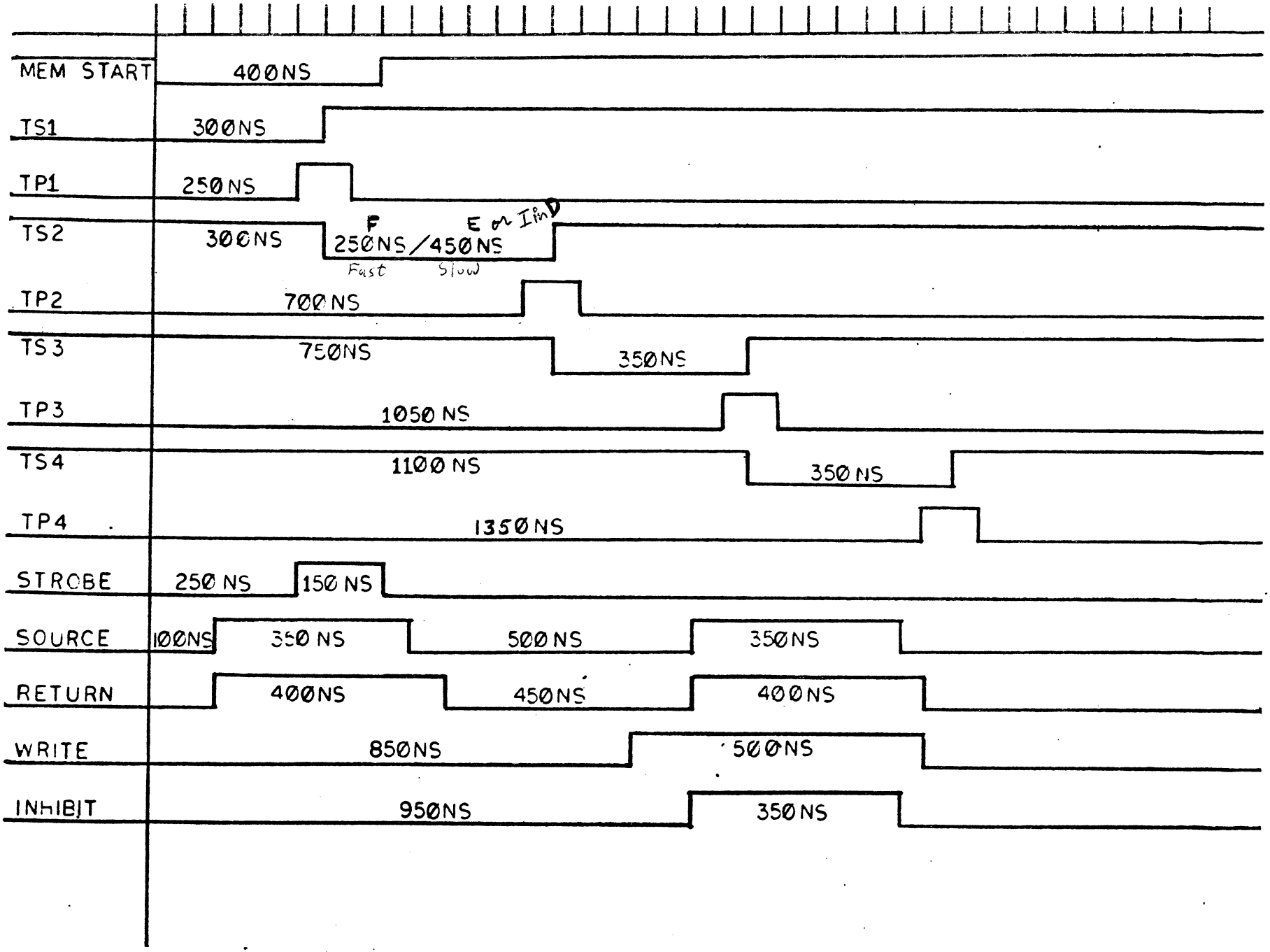
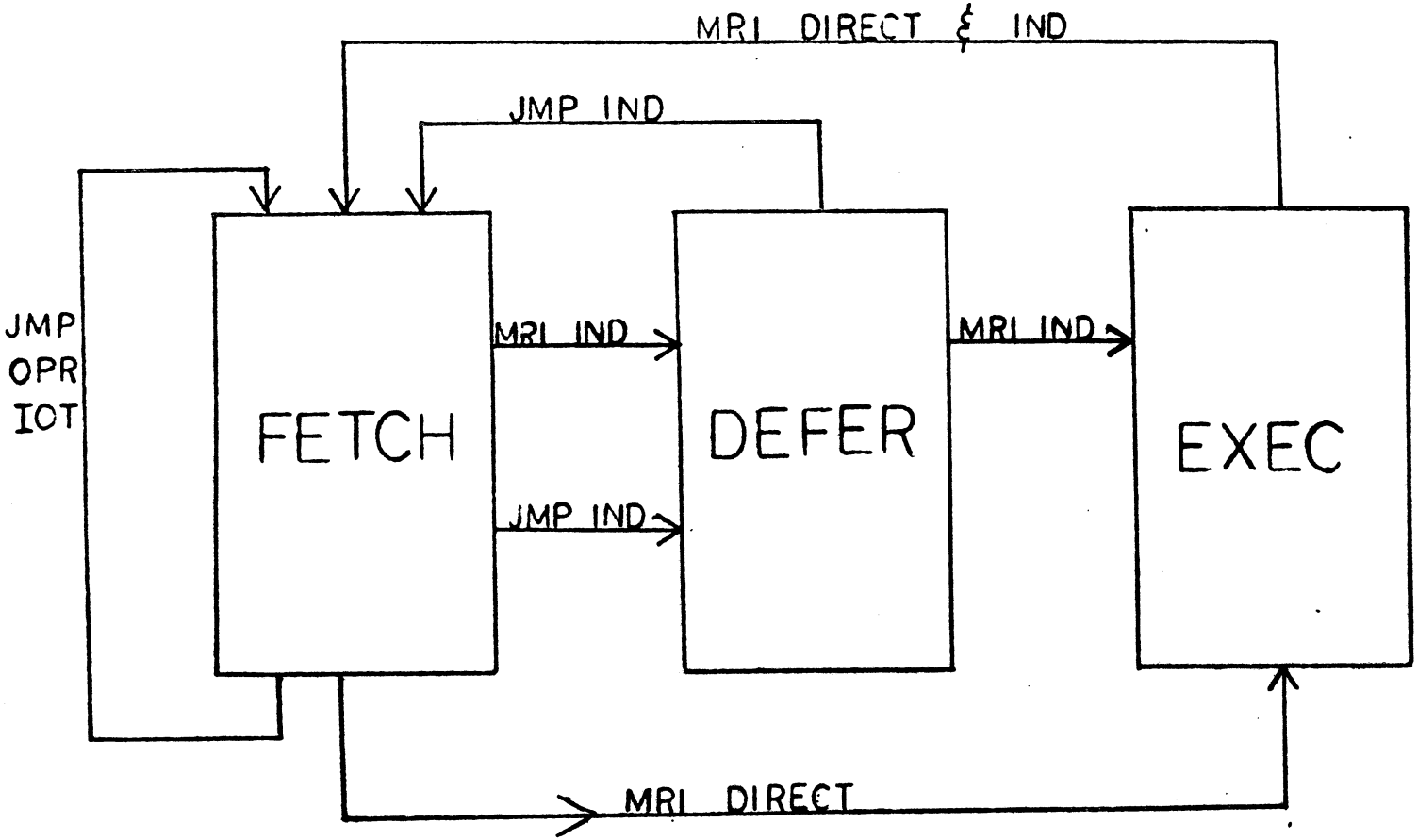
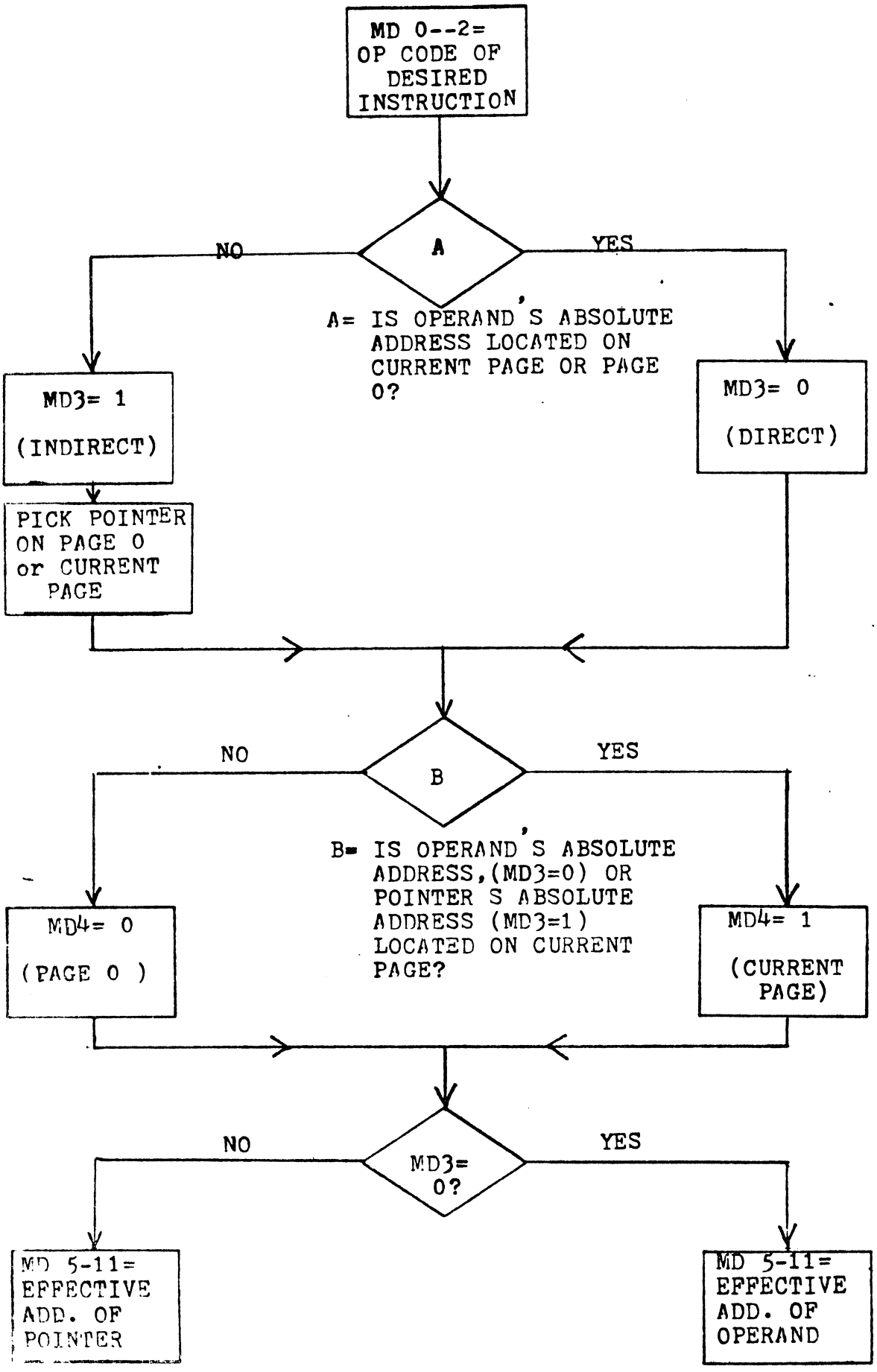


Also, See pg 3-9 in 8E Main. Man.


MEM START
400NS
TS1
300NS
TP1
250NS
TS2
300NS
F
250NS/450NS
Fast
Slow
E or IND
TP2
700NS
TS3
750NS
350NS
TP3
1050 NS
TS4
1100 NS
350 NS
TP4
1350NS
STROBE
250 NS
150 NS
SOURCE
100NS
350 NS
500NS
350NS
RETURN
400NS
450NS
400NS
WRITE
850NS
500NS
INHIBIT
950NS
350 NS

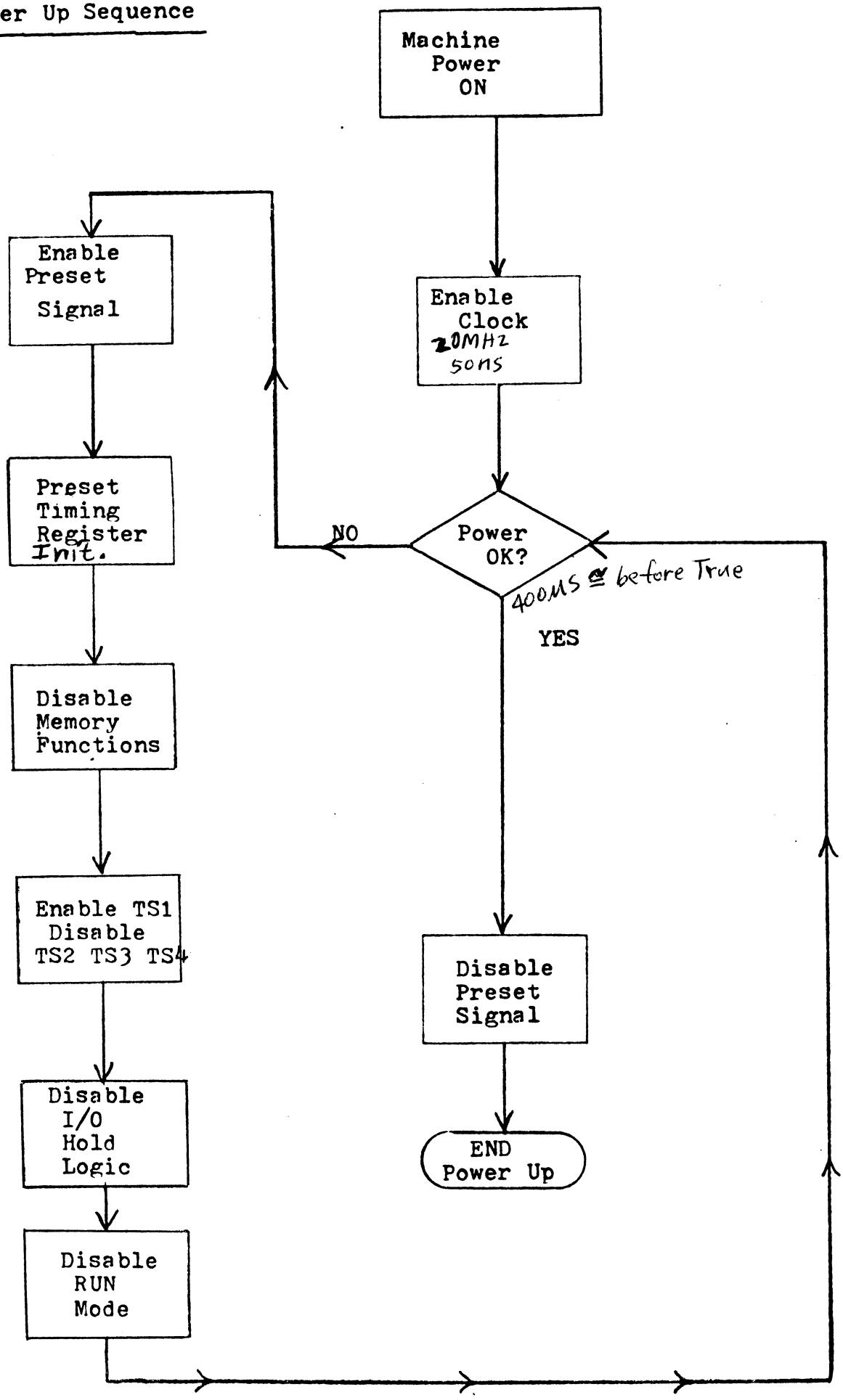

MRI DIRECT & IND
JMP IND
JMP
OPR
IOT
FETCH
MRI IND
JMP IND
DEFER
MRI IND
EXEC
MRI DIRECT

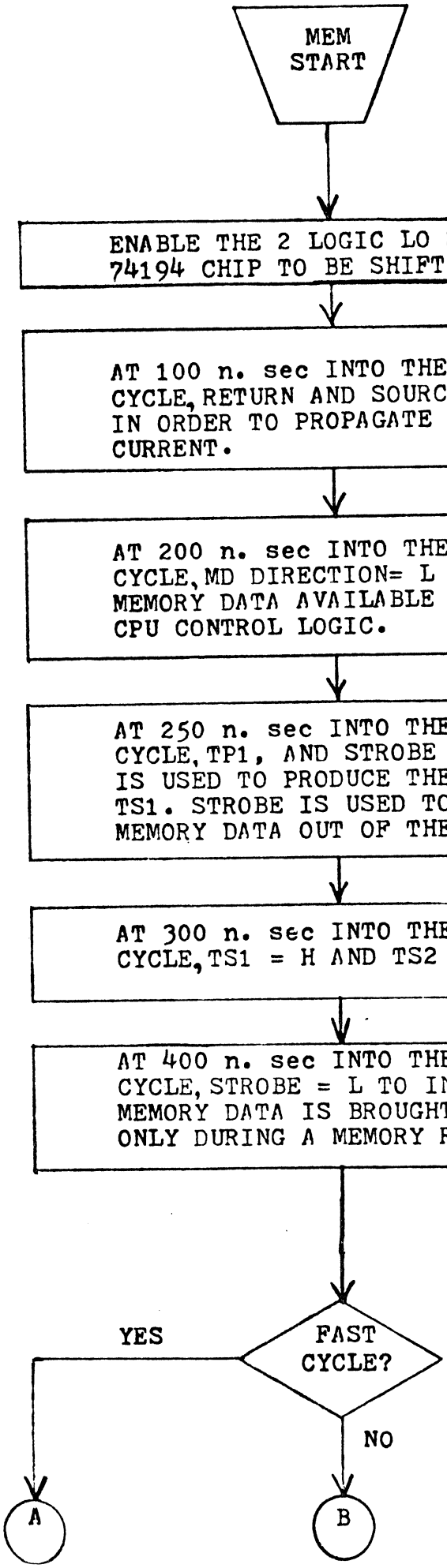
PDP 8/e
MRI INSTUCTION CODING
FLOW DIAGRAM

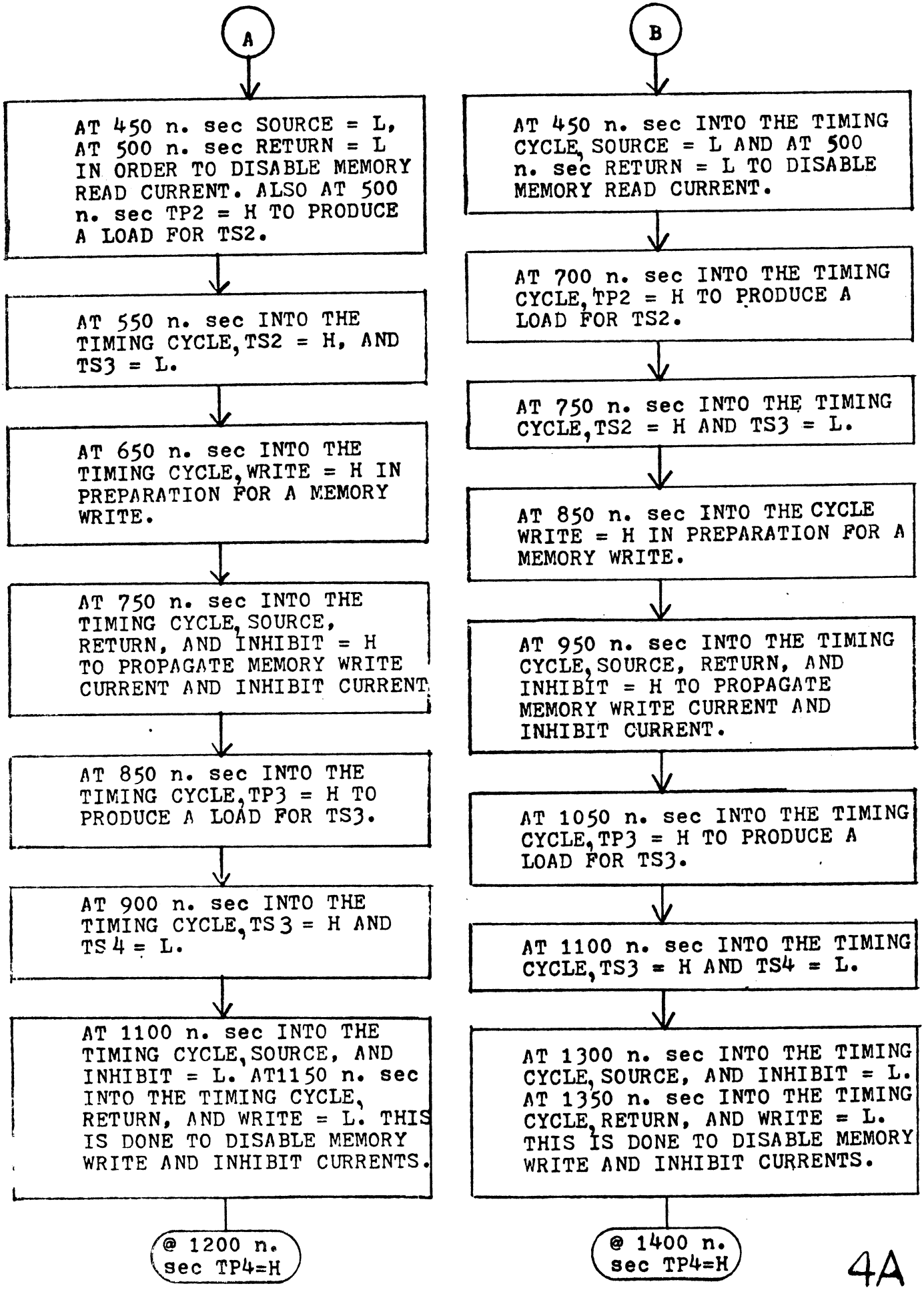


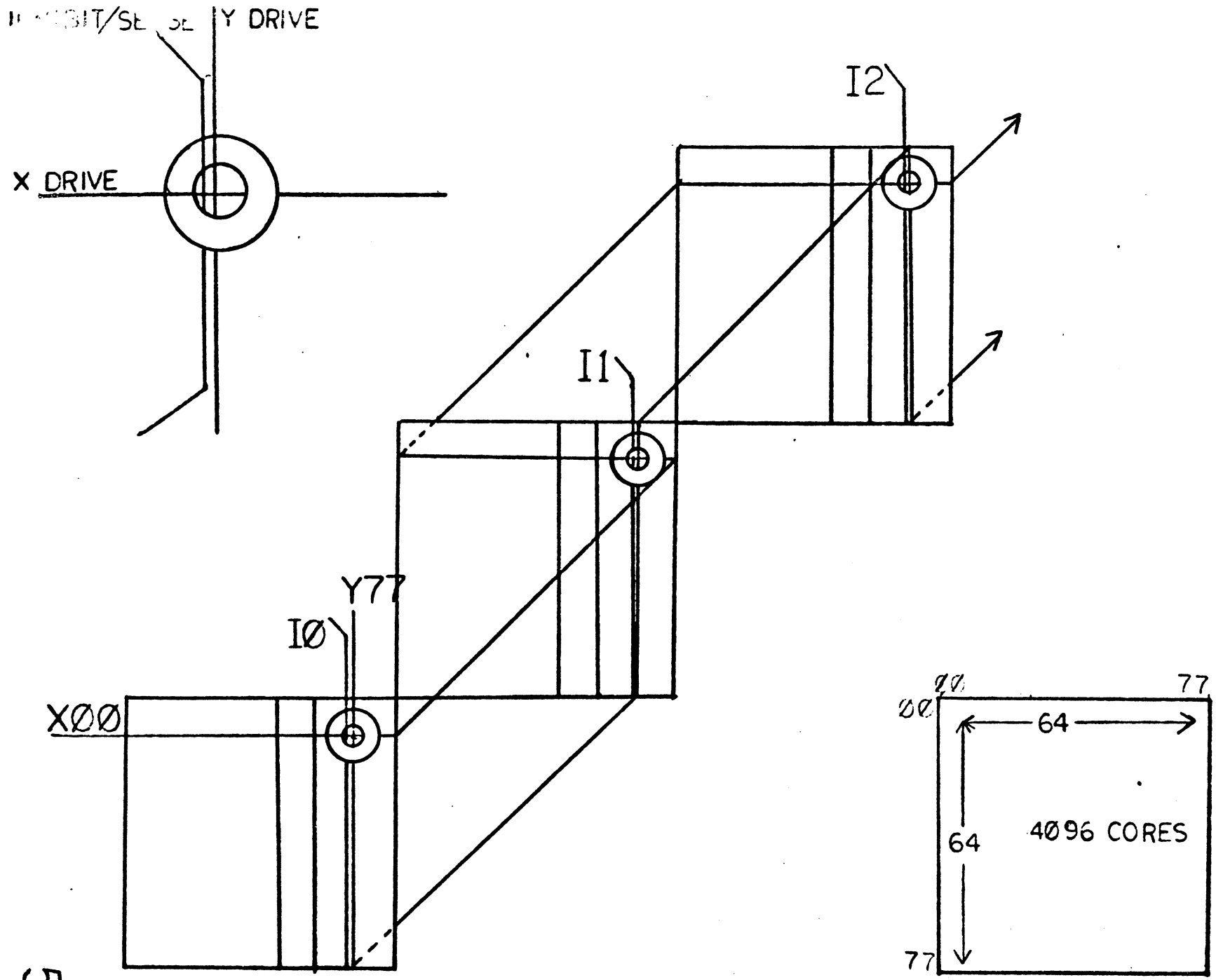
PDP 8/e

Power Up Sequence

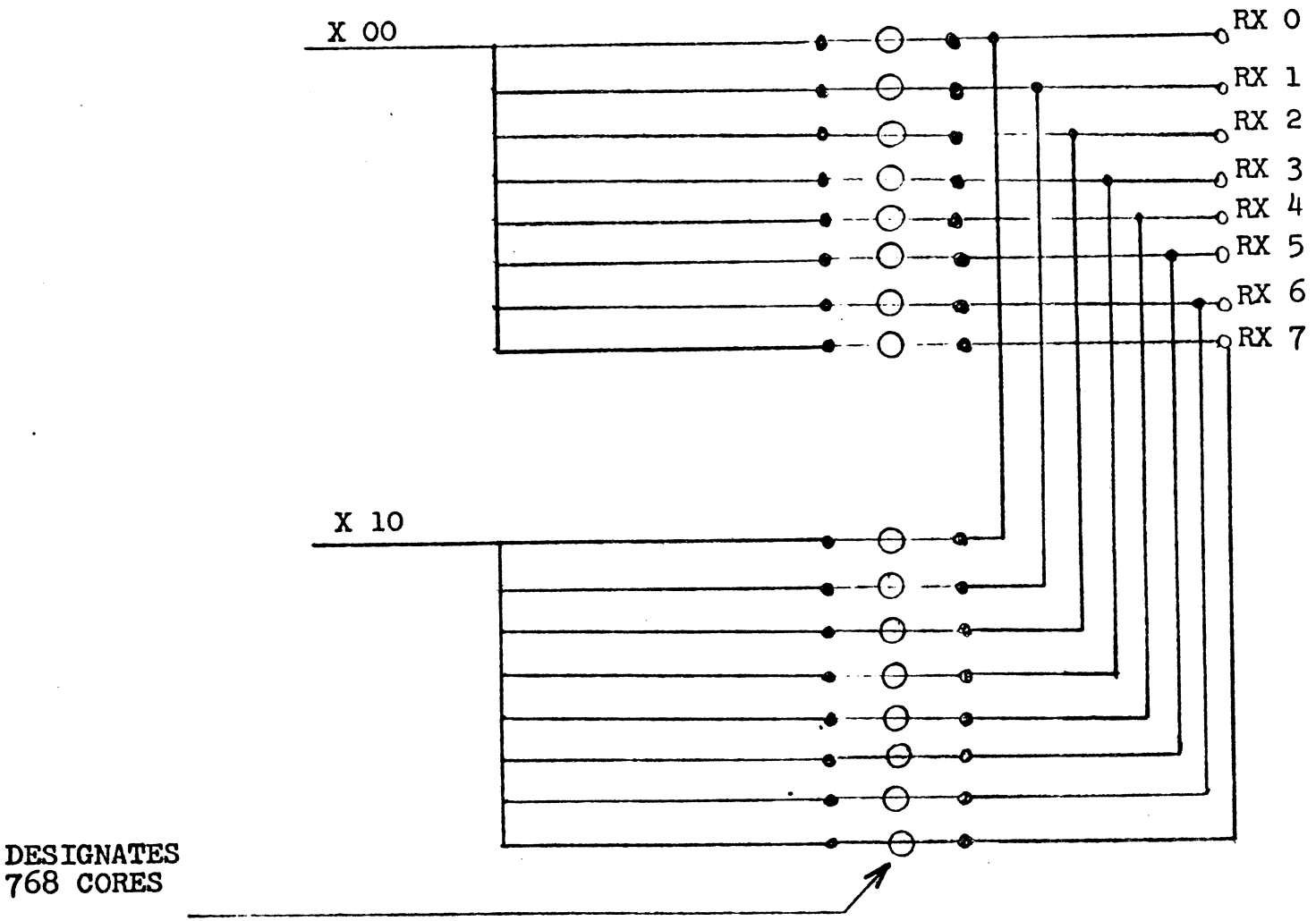



MEM START
ENABLE THE 2 LOGIC LO S OF FIRST 74194 CHIP TO BE SHIFTED RIGHT.
AT 100 n. sec INTO THE TIMING CYCLE, RETURN AND SOURCE = H IN ORDER TO PROPAGATE MEMORY READ CURRENT.
AT 200 n. sec INTO THE TIMING CYCLE, MD DIRECTION= L TO MAKE MEMORY DATA AVAILABLE TO THE CPU CONTROL LOGIC.
AT 250 n. sec INTO THE TIMING CYCLE, TP1, AND STROBE =H. TP1 IS USED TO PRODUCE THE LOAD FOR TS1. STROBE IS USED TO ENABLE THE MEMORY DATA OUT OF THE SENSE AMPS.
AT 300 n. sec INTO THE TIMING CYCLE, TS1 = H AND TS2 = L.
AT 400 n. sec INTO THE TIMING CYCLE, STROBE = L TO INSURE THAT MEMORY DATA IS BROUGHT FROM CORE ONLY DURING A MEMORY READ.
FAST CYCLE?
YES
NO
A
B


A
AT 450 n. sec SOURCE = L, AT 500 n. sec RETURN = L IN ORDER TO DISABLE MEMORY READ CURRENT. ALSO AT 500 n. sec TP2 = H TO PRODUCE A LOAD FOR TS2.
AT 550 n. sec INTO THE TIMING CYCLE, TS2 = H, AND TS3 = L.
AT 650 n. sec INTO THE TIMING CYCLE, WRITE = H IN PREPARATION FOR A MEMORY WRITE.
AT 750 n. sec INTO THE TIMING CYCLE, SOURCE, RETURN, AND INHIBIT = H TO PROPAGATE MEMORY WRITE CURRENT AND INHIBIT CURRENT.
AT 850 n. sec INTO THE TIMING CYCLE, TP3 = H TO PRODUCE A LOAD FOR TS3.
AT 900 n. sec INTO THE TIMING CYCLE, TS3 = H AND TS4 = L.
AT 1100 n. sec INTO THE TIMING CYCLE, SOURCE, AND INHIBIT = L. AT1150 n. sec INTO THE TIMING CYCLE, RETURN, AND WRITE = L. THIS IS DONE TO DISABLE MEMORY WRITE AND INHIBIT CURRENTS.
@ 1200 n. sec TP4=H
B
AT 450 n. sec INTO THE TIMING CYCLE, SOURCE = L AND AT 500 n. sec RETURN = L TO DISABLE MEMORY READ CURRENT.
AT 700 n. sec INTO THE TIMING CYCLE, TP2 = H TO PRODUCE A LOAD FOR TS2.
AT 750 n. sec INTO THE TIMING CYCLE, TS2 = H AND TS3 = L.
AT 850 n. sec INTO THE CYCLE WRITE = H IN PREPARATION FOR A MEMORY WRITE.
AT 950 n. sec INTO THE TIMING CYCLE, SOURCE, RETURN, AND INHIBIT = H TO PROPAGATE MEMORY WRITE CURRENT AND INHIBIT CURRENT.
AT 1050 n. sec INTO THE TIMING CYCLE, TP3 = H TO PRODUCE A LOAD FOR TS3.
AT 1100 n. sec INTO THE TIMING CYCLE, TS3 = H AND TS4 = L.
AT 1300 n. sec INTO THE TIMING CYCLE, SOURCE, AND INHIBIT = L. AT 1350 n. sec INTO THE TIMING CYCLE, RETURN, AND WRITE = L. THIS IS DONE TO DISABLE MEMORY WRITE AND INHIBIT CURRENTS.
@ 1400 n. sec TP4=H


Y DRIVE
X DRIVE
I2
I1
I0
Y77
X00
00
77
64
64
4096 CORES
77

EXAMPLE OF X READ LINE LAYOUT



MEMORY ADDRESS SELECTION

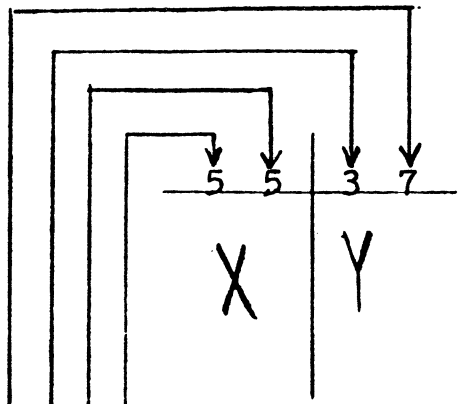


FIG. #1

THIS OCTAL DIGIT, WHICH IS COMPRISED OF MA 0---2, WILL SELECT A READ OR WRITE LINE IN THE UPPER LEFT HAND QUADRANT.

THIS OCTAL DIGIT, WHICH IS COMPRISED OF MA 3---5, WILL SELECT A READ OR WRITE LINE IN THE UPPER RIGHT HAND QUADRANT.

THIS OCTAL DIGIT, WHICH IS COMPRISED OF MA 6---8, WILL SELECT A READ OR WRITE LINE IN THE LOWER LEFT HAND QUADRANT.

THIS OCTAL DIGIT, WHICH IS COMPRISED OF MA 9---11, WILL SELECT A READ OR WRITE LINE IN THE LOWER RIGHT HAND QUADRANT.

FIG. #2

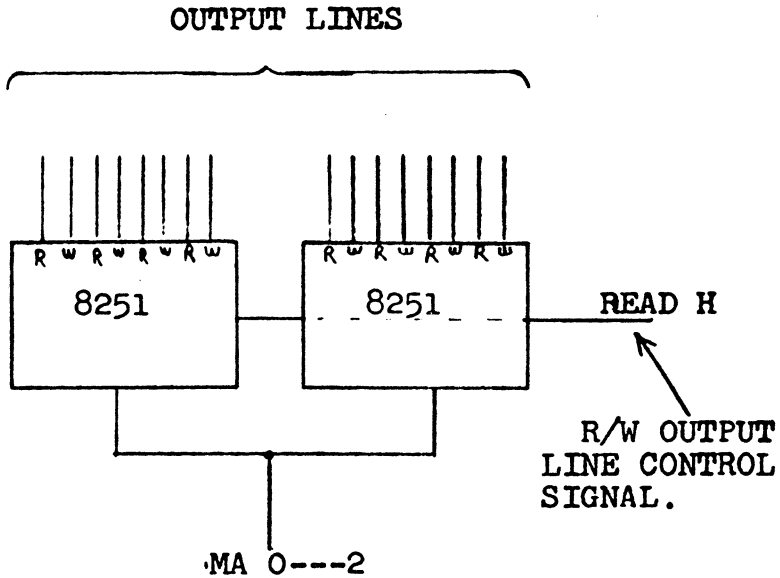
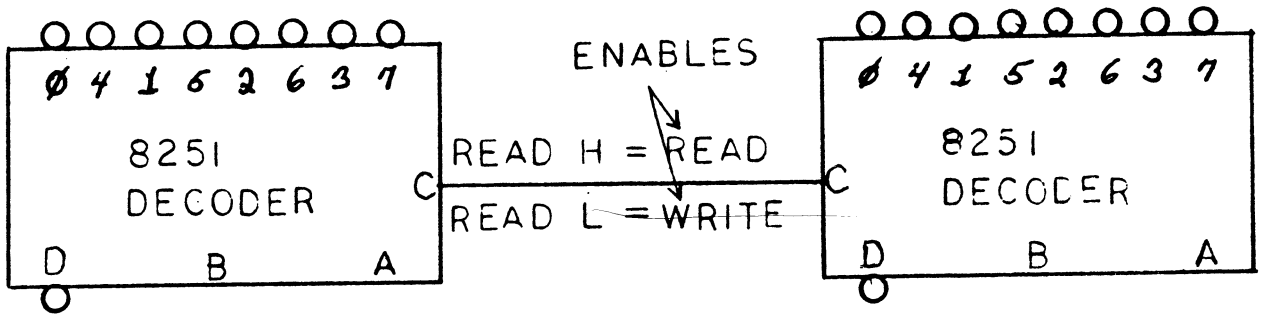


FIG. #2. ILLUSTRATES THE UPPER LEFT HAND QUADRANT OF PRINT 3. EACH 8251 HAS 8 OUTPUT LINES; 4 FOR READ AND 4 FOR WRITE FUNCTIONS. ONE READ OR WRITE LINE(DEPENDING ON THE FUNCTION) MUST BE SELECTED IN ORDER TO SELECT THE DESIRED ADDRESS. THIS IS DONE BY THE COMBINATION OF THE MA BITS(THE OUTPUT LINES WIRED AS SHOWN IN HANDOUT). THE ABOVE INFORMATION IS APPLICABLE TO THE OTHER 3 QUADRANTS OF PRINT 3.

MEMORY ADDRESS DECODER

WHEN READING, USE OUT ACTIVE PINS 4,5,6,7
WHEN WRITING, USE OUT ACTIVE PINS 0,1,2,3



MA
0 1 2
3 4 5
6 7 8
9 10 11

0	0	0
0	0	1
0	1	0
0	1	1

MA INPUT

D	B	A
L	L	L
L	L	H
L	H	L
L	H	H

OUT ACTIVE PINS

4	0
5	1
6	2
7	3

READ WRITE

MA
0 1 2
3 4 5
6 7 8
9 10 11

1	0	0
1	0	1
1	1	0
1	1	1

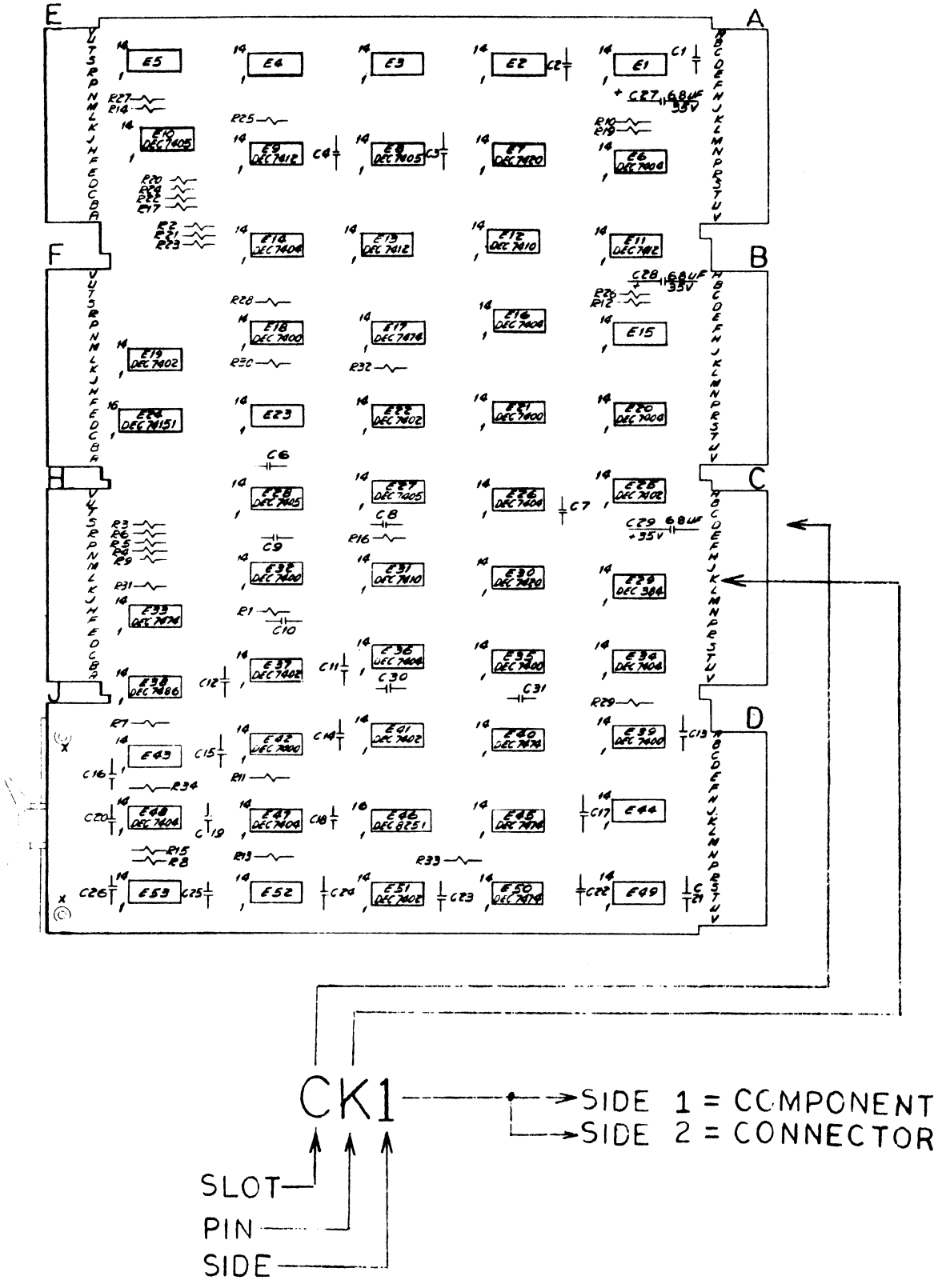
MA INPUT

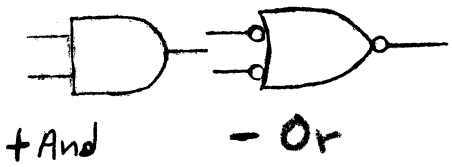
D	B	A
L	L	L
L	L	H
L	H	L
L	H	H

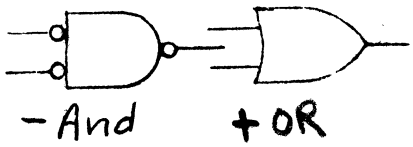
OUT ACTIV. PINS

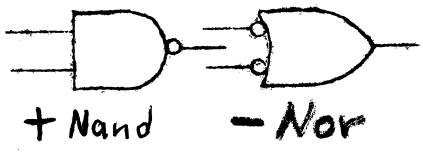
4	0
5	1
6	2
7	3

READ WRITE


CK1
SIDE 1 = COMPONENT
SIDE 2 = CONNECTOR
SLOT
PIN
SIDE



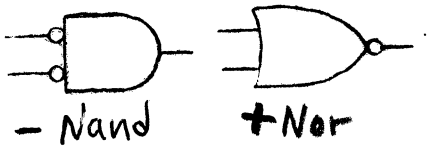




A	B	C
H	H	H
H	L	L
L	H	L
L	L	L

A	B	C
H	H	H
H	L	H
L	H	H
L	L	L

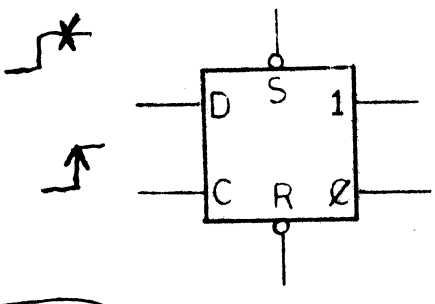
A	B	C
H	H	L
H	L	H
L	H	H
L	L	H

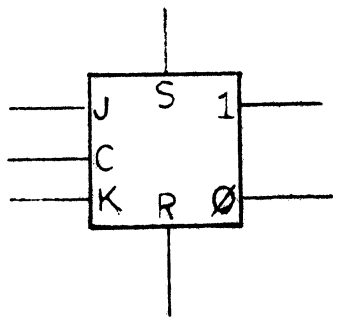


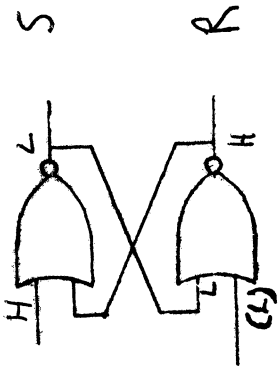
A	B	C
H	H	L
H	L	L
L	H	L
L	L	H

H	L	H
L	H	L

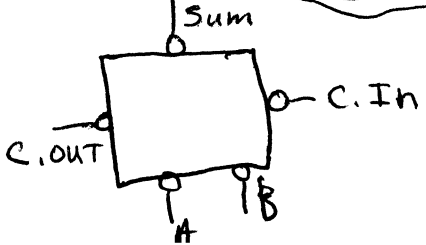
A	B	C
H	H	L
H	L	H
L	H	H
L	L	L







JK



A	B	C.I.	Sum	C.O.
1	0	0	1	0
1	1	0	0	1
1	1	1	1	1
0	0	1	1	0

L = 1
H = 0

ADDER LOGIC TRUTH TABLES

EN0	EN1	EN2	RESULT
L	L	L	PC
L	L	H	MD
L	H	L	MQ
L	H	H	MA
H	X	X	ZERO

DATA T	DATA F	RESULT
L	L	COMPLEMENT
L	H	TRUE
H	L	ZERO

PAGE	RIGHT	LEFT	TWICE	RESULT
L	L	L	L	CURRENT PAGE
X	L	L	H	AND
X	L	H	L	RIGHT 2
X	L	H	H	RIGHT 1
X	H	L	L	LEFT 2
X	H	L	H	LEFT 1
X	H	H	L	SWAP
X	H	H	H	NO SHIFT
H	L	L	L	PAGE ZERO

I/O SIGNALS

C0	C1	C2	RESULT
L	L	L	DATA ⟶ PC
L	L	H	DATA ⟶ AC
L	H	L	PC + DATA ⟶ PC
L	H	H	AC ⟶ DATA, CLR AC
H	L	L	DATA ⟶ PC
H	L	H	AC/DATA ored ⟶ AC
H	H	L	PC + DATA ⟶ PC
H	H	H	AC/DATA ored ⟶ AC

ADDER CONDITIONS

LEFT	RIGHT	CAR IN	SUM	CAR OUT
H	H	H	H	H
H	L	H	L	H
L	L	H	H	L
H	H	L	L	H
H	L	L	L	H
L	L	L	L	L

RIGHT	LEFT	TWICE	PAGE Z	USE
L	L	L	L	CURRENT PAGE
L	L	H	X	AND
L	H	L	X	RTR
L	H	H	X	RAR
H	L	L	X	RTL
H	L	H	X	RAL
H	H	L	X	BYTE SWAP
H	H	H	X	NO SHIFT
L	L	L	H	PAGE ZERO

LEFT	RIGHT	CI	SUM	CO
H	H	H	H	H
H	L	H	L	H
L	L	H	H	L
H	H	L	L	H
H	L	L	H	L
L	L	L	H	L

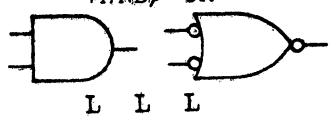
ENO	EN1	EN2	ADDER IN
L	L	L	PC
L	L	H	MD
L	H	L	MQ
L	H	H	MA
H	X	X	ZERO

DATA T	DATA F	ADDER IN
L	L	DATA BUS
L	H	DATA BUS
H	L	ZERO

IND1	IND2	RESULT
L	L	AC → BUS
L	H	BUS
H	L	MQ → BUS
H	H	STATUS

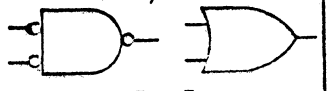
C0	C1	C2	RESULT
L	L	L	DATA → PC
L	L	H	DATA → AC
L	H	L	PC + DATA → PC
L	H	H	AC → DATA, 0 → AC
H	L	L	DATA → PC
H	L	H	AC ored DATA → AC
H	H	L	PC + DATA → PC
H	H	H	AC ored DATA → AC

+AND/-OR



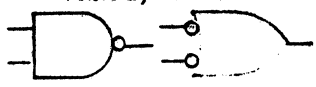
L	L	L
L	H	L
H	L	L
H	H	H

-AND/+OR



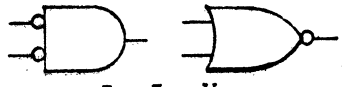
L	L	L
L	H	H
H	L	H
H	H	H

+NAND/-NOR



L	L	H
L	H	H
H	L	H
H	H	L

-NAND/+NOR

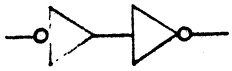


L	L	H
L	H	L
H	L	L
H	H	L

XOR

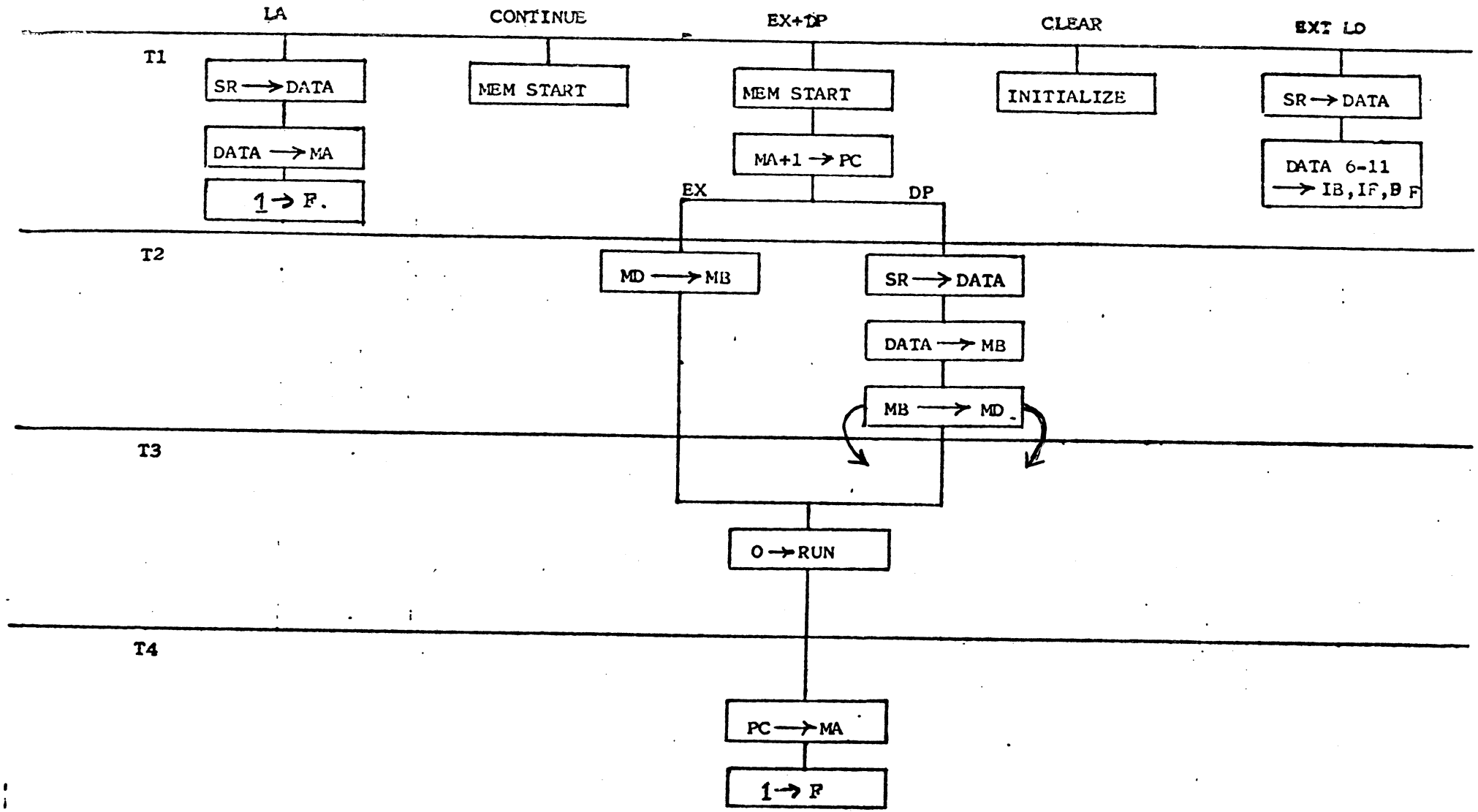
L	L	L
L	H	H
H	L	H
H	H	L

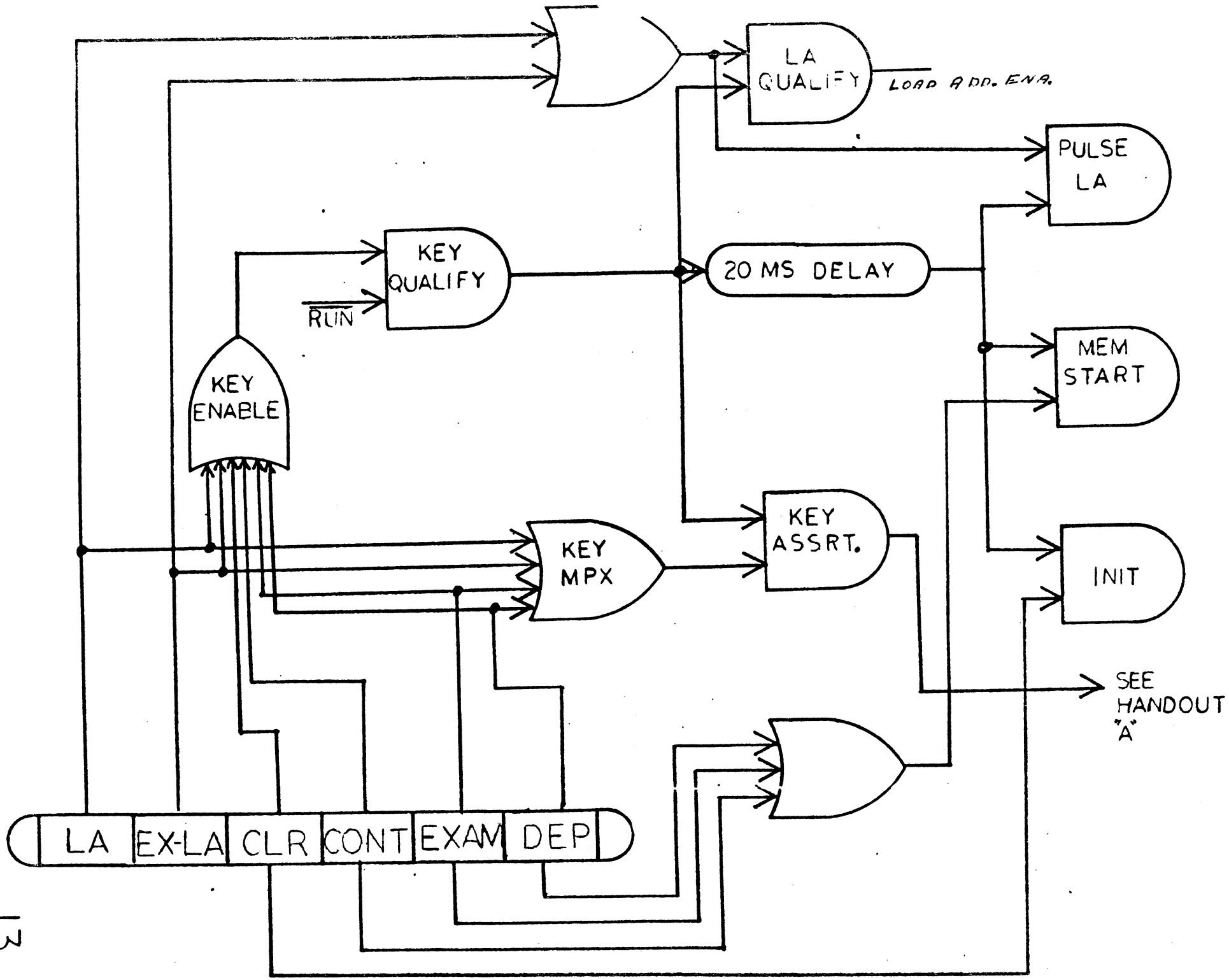
INVERTER

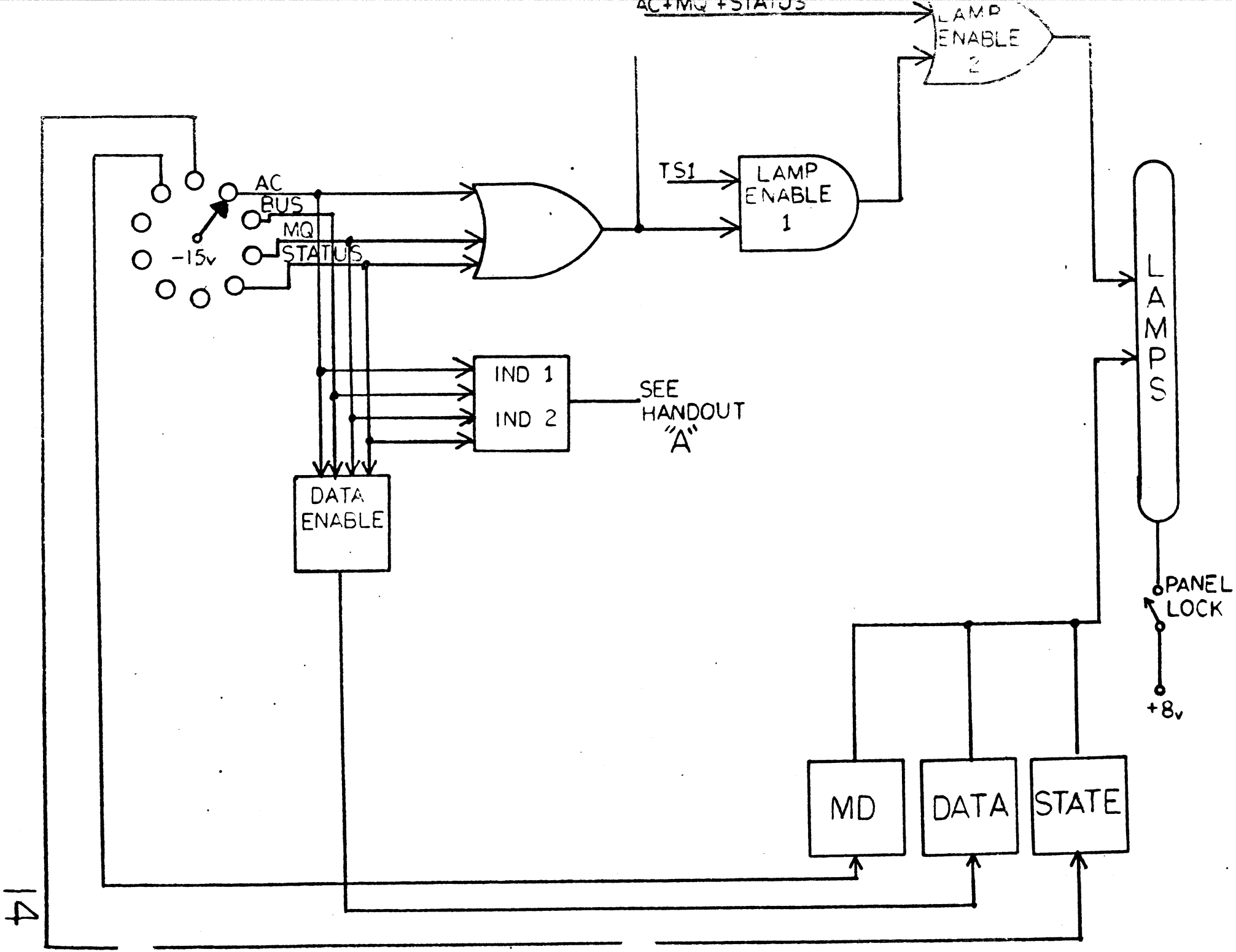


L	H	L
H	L	H

PDP - 8E MANUAL KEYS FLOW DIAGRAM




LA QUALIFY
LOAD ADD. ENA.
PULSE LA
KEY QUALIFY
RUN
20 MS DELAY
MEM START
KEY ENABLE
KEY MPX
KEY ASSRT.
INIT
SEE HANDOUT "A"
LA
EX-LA
CLR
CONT
EXAM
DEP


AC+MQ +STATUS
LAMP ENABLE 2
TS1
LAMP ENABLE 1
AC BUS
MQ
STATUS
-15v
IND 1
IND 2
SEE HANDOUT "A"
DATA ENABLE
LAMPS
PANEL LOCK
+8v
MD
DATA
STATE

PDP 8/e

KEY FUNCTION SIGNALS

HANDOUT A

IND 1	IND 2.	Result
L	L	AC ⟶ BUS
L	H	BUS ⟶ BUS
H	L	MQ ⟶ BUS
H	H	STATUS ⟶ BUS

KEY LA

(A)	KEY CONTROL	H	Qualification for MA LOAD
(B)	MS DISABLE	L	Disables MAJOR STATES and IR
(C)	MD DIR. IN.	L	NO use for KEY LA
(D)	BK DATA CONTROL	L	NO use for KEY LA
(E)	SR----BUS	H	Enables SR to the DATA BUS

KEY EXAM

(A)	KEY CONTROL	L	Allows one timing cycle only
(B)	MS DISABLE	L	Disables MAJOR STATES and IR
(C)	MD DIR, IN.	L	Enables MEM ⟶ MD LINES
(D)	BK DATA CONTROL	L	Enables MD ⟶ MB
(E)	SR----BUS	L	Disables SR ⟶ BUS

KEY DEP

(A)	KEY CONTROL	L	Allows one timing cycle only
(B)	MS DISABLE	L	Disables MAJOR STATES and IR
(C)	MD DIR. IN.	H	Disables MEM ⟶ MD LINES
(D)	BK DATA CONTROL	H	Disables MD ⟶ MB
(E)	SR----BUS	H	Enables SR ⟶ BUS

FETCH FOR AND, TAD, ISZ, DCA, JMS, JMP I

MS FLOW SHEET 1

FUNCTION	EN0	EN1	EN2	DATA T	DATA F	CARRY IN	MD DIR	RIGHT	LEFT	TWICE	PAGE ZERO	TIME
MA+1 → PC	L	H	H	H	L	L	L	H	H	H	L	TS1
AC → [illegible] DATA MD → MB MD 0-2 → IR	L	L	H	H	L	H	L	H	H	H	L	TS2
↓	H	H	H	H	L	H	L	H	H	H	L	TS3
MD 5-11 → CPMA (0) (1) 0 → MA 0-4 MA → MA 0-4	H	H	H	H	L	H	L	L	L	L	MD4L L - - - - - MD4L H	TS4

FUNCTION	EN0	EN1	EN2	DATA T	DATA F	CARRY IN	MD DIR	RIGHT	LEFT	TWICE	PAGE Z	TIM
AND,TAD,ISZ, DCA,JMSP0	H	H	H	H	L	H	L	H	H	H	L	TS1
AND,TAD	L	L	H	H	L	H	L	H	H	H	L	TS2
ISZ	L	L	H	H	L	L	L	H	H	H	L	TS2
DCA	H	H	H	L	H	H	L	H	H	H	L	TS2
JMS	L	L	L	H	L	H	L	H	H	H	L	TS2
AND	H	H	H	H	L	H	H	L	L	H	L	TS3
TAD	L	L	H	L	H	H	H	H	H	H	L	TS3
ISZ	H	H	H	H	L	H	H	H	H	H	L	TS3
DCA	H	H	H	H	L	H	H	H	H	H	L	TS3
JMS	L	H	H	H	L	L	H	H	H	H	L	TS3
AND,TAD,ISZ, DCA,JMS	L	L	L	H	L	H	H	H	H	H	L	TS4

FUNCTION	EN0	EN1	EN2	DATA T	DATA F	CARRY IN	MD DIR	RIGHT	LEFT	TWICE	PAGE Z	TIM
↓	H	X	X	H	L	H	L	H	H	H	L	TS
MD + 1 → MB	L	L	H	H	L	L	L	H	H	H	L	TS
JMP MD → PC	L	L	H	H	L	H	L	H	H	H	L	TS
$\overline{JMP}$	H	X	X	H	L	H	AUTO INDEX H	H	H	H	L	
MD → CPMA	L	L	H	H	L	H	L	H	H	H	L	TS
JMP PC → CPMA	L	L	L	H	L	H	AUTO INDEX H	H	H	H	L	

PDP-8E MECHANIZATION TABLE

FETCH FOR AND, TAD, ISZ, DCA, JMS INSTRUCTIONS

Function	Source Req'd	Route Req'd	Dest'n Req'd	Enabling Levels & Boolean Expression	Source Page & Cord	Dest'n Page & Cord	Time	Major State
MA+1→PC	CPMA			ENOL, EN1H, EN2H = TS1 (BF + KC) "Key Cont"	11H7	6C8	TS1	FETCH
		+1		CARRY IN L = TS1 (BF + KC)	11H1	9		
			PC	PC LOAD = TP1 (BF + KC)	10J5	6B2	TP1	
0→SKIP				0→SKIP = BF · TP1	10H1	9J1		
SA→MD	SA			TIME STROBE = (STROBE · FIELD · $\overline{\text{WRITE}}$) DELAYED	5B8	5E1-8	TS2	
			MD LINES	MD DIR L = 100NS INTO TS1, 0→MD DIR	14C7	5E1-8		
AC→DATA	AC			AC→BUS L = BF · TS2	11J3	6C6		
MD→MB	MD LINES			ENOL, EN1L, EN2H = ($\overline{\text{DCA}}$ · E · TS2) · BTS2 · ($\overline{\text{MS DIS}}$ + BRK DATA CONT)	11H5	6C8		
			MB	MB LOAD = BTP2	10H2	6C2	TP2	
MD0-2→IR			IR	LOAD IR = TP2 · F	12E5	12E5		
							TS3	
							TP3	
MD 5-11 →CPMA	MD5-11	FOR PAGING		RIGHT L, LEFT L, TWICE L = (BTS4 · $\overline{\text{F SET}}$ · BF · $\overline{\text{FE + FD}}$)	10J5	6F8	TS4	
MD 4(0) 0→CPMA 0-4	0			PAGE 2 H = (MD4(0) · (BTS4 · $\overline{\text{FSET}}$ · BF · $\overline{\text{FE + FD}}$)	10J5	6F8		
MD4(1) CPMA 0-4 →CPMA 0-4	CPMA 0-4			PAGE 2 L = MD4(1) · (BTS4 · $\overline{\text{FSET}}$ · BF · $\overline{\text{FE + FD}}$)	10J5	6F8		
			CPMA	CPMA LOAD = TP4 · $\overline{\text{MALC}}$	12H4	6F3	TP4	
MD 3(0) I→E			EXECUTE	CPMA LOAD (F · $\overline{\text{OPR + IOT}}$ · MD3(0) · $\overline{\text{JMP}}$) = E L	12B7	12B4		
MD 3(1) I→D			DEFER	CPMA LOAD (F · $\overline{\text{OPR+IOT}}$ · MD3(1)) = D L	12B7	12B4		

PDP-8E MECHANIZATION TABLE

DEFER FOR AND, TAD, ISZ, DCA, JMS INSTRUCTIONS

Function	Source Req'd	Route Req'd	Dest'n Req'd	Enabling Levels & Boolean Expression	Source Page & Cord	Dest'n Page & Cord	Time	Major State	Instruction
							TS1	DEFER	
							TP1		
SA→MD	SA			TIME STROBE = (STROBE · FIELD · $\overline{\text{WRITE}}$) DELAYED	5B8	5E1-8	TS2		
			MD LINES	MD DIR L = 100NS INTO TS1, 0→MD DIR	14C7	5E1-8			
MD+1→MB	MD			ENOL, EN1L, EN2H = ($\overline{\text{DCA}}$ · E · $\overline{\text{TS2}}$) · BTS2 · ($\overline{\text{MS DIS}}$ + BRK DATA CONT)	11H 5	6C8			
		+1		CARRY IN L = BD · BTS2	11H1	9			
			MB	MB LOAD = BTP2	10H2	6 C2	TP2		
AUTO INDEX MB→MD MD DIR H	MB			MD DIR H = TP2 (D · AUTO INDEX)	14C7	6 C3			
AUTO INDEX MEM→MD MD DIR L	MEM			MD DIR L = TP2 (D · $\overline{\text{AUTO INDEX}}$)	14C7	5E8			
							TS3		
							TP3		
MD→CPMA	MD			ENOL, EN1L, EN2H = $\overline{\text{BTS4}}$ · (D · $\overline{\text{JMP}}$)	11H5	6 C8	TS4		
			CPMA	CPMA LOAD = TP4 · $\overline{\text{MALC}}$	10H4	6 F3	TP4		
1→E			EXECUTE	CPMA LOAD · (D · $\overline{\text{JMP}}$) = E L	12B7	12B4			

PDP-8L MECHANIZATION TABLE

EXECUTE FOR AND, TAD, ISZ, JMS INSTRUCTIONS

Function	Source Req'd	Route Req'd	Dest'n Req'd	Enabling Levels & Boolean Expression	Source Page & Cord	Dest'n Page & Cord	Time	Major State	Instruction
							TS1	Execute	ALL
							TP1		
SA→MD	SA			TIME STROBE = (STROBE • FIELD • WRITE)	5B8	5E1-8	TS2		ALL
			MD LINES	MD DIR L = 100 NS INTO TSI, 0→MD DIR	14C7	5E1-8			
MD→MB	MD			ENOL, EN1L, EN2H = (DCA•E•TS2) BTS2• (MS DIS + BRK DATA CONT)	11H5	6C8			AND TAD
MD+1→MB	MD			ENOL, EN1L, EN2H = (DCA•E•TS2) BTS2• (MS DIS + BRK DATA CONT)	11H5	6C8			ISZ
		+1		CARRY IN = ISZ • E • TS2	11H2	9			
AC→DATA	AC			AC→BUS = DCA • E • TS2	11J3	6C6			DCA
DATA→MB		BUS ENABLE		DATA L, DATA F H = DCA • E • TS2	11J7	6D6			
				ENOL, EN1L, EN2L = JMS • E • TS2	11H6	6C8			JMS
PC→MB	PC				10H2	6C2	TP2		ALL
			MB	MB LOAD = BTP2	14C7	6 C3			ALL
MB→MD	MD DIRH			MD DIR H = F + (D • AUTO INDEX)					
CARRY OUT →OVERFLOW				BTP3 • CARRY OUT = OVERFLOW	10D7	10H3			
AC∧MB→ AC	AC∧MB			LEFT L, RIGHT L, TWICE H = AND • E • TS3	10H6	6 F8	TS3		AND
AC + MD→ AC	AC			AC→BUS = TAD • E • TS3	11J3	6 C6			TAD
		BUS ENABLE		DATA T L, DATA F H = TAD • E • TS3	11J7	6D6			
	MD			ENO L, EN1L, EN2H = TAD • E • TS3	11H5	6C8			
				ENO H, EN1H, EN2H	11J5	6C8			DCA
0→AC	0			ENO L, EN1H, EN2H = JMS • E • TS3	11F6	6C8			JMS
MA+1→PC	MA								
		+1		CARRY IN = (JMS • E • TS3) MA DIS • ROM ADDRESS	11H2	9			
CARRY OUT LINK→LINK			LINK	CARRY OUT • OP1 • OPE+IOT • TAD • E • BTP3 • LINK(1) = LINK (0) CARRY OUT • OP1 • OPE+IOT • TAD • E • BTP3 • LINK (0) = LINK (1)	11E-J7	11J7	TP3		TAD
OVERFLOW → SKIP				OVERFLOW • (ISZ • E • TS3) • BTP3 = 1→SKIP	10H3	9J1			ISZ
			AC	AC LOAD = BTP3 (AND•E + TAD • E + DCA • E + OPR • F)	10H3	6B2			AND TA DCA
			PC	PC LOAD = JMS • E • BTP3	10J4	6B2			JMS

$\overline{\text{SKIP}}$ PC→CPMA	PC			ENOL, EN1L, EN2L = (TS4 • F SET • $\overline{\text{MS DIS}}$) • ($\overline{\text{FE+FD}}$)	11F6	6C8	TS4		ALL
SKIP PC+1→ CPMA	PC			ENOL, EN1L, EN2L = (TS4 • F SET • $\overline{\text{MS DIS}}$) • ($\overline{\text{FE+FD}}$)	11F6	6C8			ISZ
		+1		CARRY IN = SKIP (1) • (TS4 • F SET $\overline{\text{MS DIS}}$	11H1	9			
			CPMA	TP4 • $\overline{\text{MALC}}$ = CPMA LOAD	10H4	6F3	TP4		ALL
1→F			FETCH	($\overline{\text{F+D+FE SET + F D SET}}$) = F SET F SET • CPMA LOAD = FL	12B4	12B4			

PDP-8E MECHANIZATION TABLE
FETCH FOR JMP AND JMPI INSTRUCTIONS

Function	Source Req'd	Route Req'd	Dest'n Req'd	Enabling Levels & Boolean Expression	Source Page & Cord	Dest'n Page & Cord	Time	Major State	Instruction
MA+1→PC	CPMA			ENOL, EN1H, EN2H = TS1 (BF + KC)	11H7	6C8	TS1	FETCH	ALL
		+1		CARRY IN L = TS1 (BF + KC)	11H1	9			
			PC	PC LOAD = TP1 (BF + KC)	10J5	6B2	TP1		
				0→SKIP = BF · TP1	10H1	9J1			
0→SKIP				TIME STROBE = (STROBE · FIELD · WRITE) DELAYED	5B8	5E1-8	TS2		
SA→MD	SA								
			MD LINES	MD DIR L = 100NS INTO TS1, 0→MD DIR	14C7	5E1-8			
				AC→BUS L = BF · TS2	11J3	6C6			
AC→DATA	AC			ENOL, EN1L, EN2H = (DCA·E·TS2)·BTS2·(MS DIS + BRK DATA CONT)	11H5	6C8			
MD→MB	MD LINES								
			MB	MB load = BTP2	10H2	6C2	TP2		
			IR	LOAD IR = TP2·F	12E5	12E5			
MD0-2→ IR									
MD5-11→ PC	MD5-11	PAGE		LEFT L, RIGHT L, TWICE L = BTS3 · JMP·BF	10H6	13J2	TS3		
				PAGE 2 H = MD4(0)·(BTS3·JMP·BF)	10H6	6E8			
MD 4(0) 0→PC0-4	0			PAGE 2 L = MD4(1)(BTS3·JMP·BF)	10H6	6E8			
MD 4(1) MA0-4→ PC0-4	CPMA 0-4								
			PC	PC LOAD = BTP3 · JMP	10H5	6B2	TP3		
MD 3(0) PC→MA	PC			ENOL, EN1L, EN2L = (TS4·F SET· MS DIS) ·(FE + FD)	11H6	6C8	TS4		JMP DIREC
MD 3(1) MD 5-11→ MA	MD5-11	PAGE		RIGHT L, LEFT L, TWICE L = (BTS4·F SET ·BF·FE + FD)	10J5	13J2			JMP I DIREC
MD 4(0) 0→CPMA 0-4	0			PAGE 2H = MD 4(0)·(BTS4·F SET·BF· FE + FD)	10J5	6F8			
MD 4(1) CPMA 0-4→ CPMA 0-4	CPMA 0-4			PAGE 2 L = MD4(1)·(BTS4·F SET·BF· FE + FD)	10J5	6F8			
			CPMA	CPMA LOAD = TP4 · MALC	10H4	6F3	TP4		ALL
			FETCH	CPMA LOAD (F·OPR + IOT·MD 3(0)·JMP =FL	12 B7	12B4			JMP DIREC
MD 3(0) 1→F			DEFER	CPMA LOAD (F·OPR+IOT·MD 3(1) = DL	12B7	12B4			JMP I DIREC
MD 3(1) 1→ D									

PDP-8E MECHANIZATION TABLE
DEFER FOR JMP I INSTRUCTION

Function	Source Req'd	Route Req'd	Dest'n Req'd	Enabling Levels & Boolean Expression	Source Page & Code	Destination Page & Code	Time	Major State
							TS1	DEFER
							TP1	
SA→MD	SA			TIME STROBE = (STROBE·FIELD·$\overline{\text{WRITE}}$) DELAYED	5B8	5E 1-8	TS2	
			MD LINES	MD DIR L = 100NS INTO TS1, 0→MD DIR	14C7	5E 1-8		
MD+1→ MB	MD			ENOL, EN1L, EN2H = ($\overline{\text{DCA·E·TS2}}$)·BTS2· ($\overline{\text{MS DIS}}$ + BRK DATA CONT)	11H5	6C8		
		+1		CARRY IN = BD·BTS2	11H1	9		
			MB	MB LOAD = BTP2	10H2	6C2	TP2	
AUTO INDEX MB→MD MD DIR H	MB			MB DIR H = TP2 (D · AUTO INDEX)	14C7	6C3		
$\overline{\text{AUTO INDEX}}$ MEM→MD MD DIR L	MEM			MD DIR L = TP2 (D · $\overline{\text{AUTO INDEX}}$)	14C7	5E8		
MD→PC	MD			ENOL, EN1L, EN2H = JMP · BTS3 · BD	11F4	6C8	TS3	
			PC	PC LOAD = BTP3 · JMP	10H5	6B2	TP3	
PC→CPMA	PC			ENOL, EN1L, EN2L = (TS4 · F SET · $\overline{\text{MS DIS}}$) · ($\overline{\text{FE + FD}}$)	10H6	6C8	TS4	
			CPMA	CPMA LOAD = TP4 · $\overline{\text{MALC}}$	10H4	6F3	TP4	
1→F			FETCH	JMP · D · CPMA LOAD = FL	12B7	12B4		

PDP-8E MECHANIZATION TABLE
FETCH FOR OPERATE GROUP ONE INSTRUCTIONS

Function	Source Req'd	Route Req'd	Dest'n Req'd	Enabling Levels & Boolean Expression	Source Page & Cord	Dest'n Page & Cord	Time	Major State	Instruction
MA+1→PC	CPMA			ENOL, EN1H, EN2H = TS1 (BF + KC)	11H7	6C8	TS1	FETCH	ALL
		+1		CARRY IN L = TS1 (BF + KC)	11H1	9			
			PC	PC LOAD = TP1 (BF + KC)	10J5	6B2	TP1		
0→SKIP				0→SKIP = BF • TP1	10H1	10J1			
SA→MD	SA			TIME STROBE = (STROBE•FIELD•WRITE) DELAYED	5B8	5E 1-8	TS2		
			MD LINES	MD DIR L = 100NS INTO TS1, 0→MD DIR	14C7	5E 1-8			
AC→DATA	AC			AC→BUS L = BF • TS2	11J3	6C6			
MD→MB	MD LINES			ENOL, EN1L, EN2H = (DCA•E•TS2)• BTS2•(MS DIS + BK DATA CONT)	11H5	6C8			
			MB	MB LOAD = BTP2	10H2	6C2	TP2		
MD0-2→ IR			IR	LOAD IR = TP2 • F	12E5	12 E5			
MD4(0) AC→DATA	AC			AC→BUS = (BTS3•OPR•F•MD4(0)•(OPE•TS3• MD7(1)•AC→BUS INH	11J3	6C6	TS3		NOP
MD5(1) 0→LINK	0			LINK DATA INPUT L = (MD5(1)•OP1)• (OP1•MD7(0))•CARRY OUT	10D-J7	10H7			CLL
MD6(1) DATA→AC	DATA			DATA T L, DATA F L = (OPR+IOT•C2L + C1L+C0H•BTS3)•(OP1•TS3•MD6(1))	10H7	6D6			CMA
MD7(1) LINK → LINK	0→LINK			LINK DATA INPUT L = (MD 7(1) • OP1)• (MD5(0) • OP1 • LINK(1)) • CARRY OUT	10D-J7	10H7			CML
	1→LINK			LINK DATA INPUT H = (MD 7(1) •OP1)• (MD5(0) •OP1•LINK(0) • CARRY OUT)					
MD 11(1) CARRY IN		+1		CARRY IN = OP1•TS3•MD 11(1)	11H1	9			IAC
CARRY OUT LINK→LINK	0→LINK			LINK DATA INPUT L = (OP1•MD7(0)•(OP1 •MD5(0))•LINK(1)•CARRY OUT	10D-J7	10H7			
	1→LINK			LINKDATA INPUT H = (OP1•MD7(0)•OP1• MD5(0)•LINK(0)•CARRY OUT)	10D-J7	10H7			
MD10(0)• MD8(1)• MD9(0) RAR		RIGHT 1		LEFT H, RIGHT L, TWICE H = MD8(1)• OP1•TS3	10H7	6F8			RAR
MD10(0)• MD8(0)• MD9(1) RAL		LEFT 1		LEFT L, RIGHT H, TWICE H = MD9(1)• OP1•TS3	10H7	6F8			RAL

MD 10(1)• MD 8(1)• MD 9(0) RTR		RIGHT 2		LEFT H, RIGHT L, TWICE L = (MD10(1)• OP1•TS3)•(MD 8(1)•OP1•TS3)	10H7	6F8		RTR
MD 10(1)• MD 8(0)• MD 9(1) RTL		LEFT 2		LEFT L, RIGHT H, TWICE L = (MD10(1)• OP1•RS3)•(MD9(1) •OP1•TS3)	10H7	6F8		RTL
MD 10(1)• MD 8(0)• MD 9(0) ByTE SWAP AC0-5→ AC6-11 AC6-11→ AC0-5		SWAP		LEFT H, RIGHT H, TWICE L = (MD10(1)• OP1•TS3)	10H7	6F8		ByTE SWA
DATA→AC			AC	AC LOAD = BTP3•(OPR•F)	10H3	6B2	TP3	ALL
			LINK	LINK CLOCK = BTP3•OP1	10H8	10H8		
PC→CPMA	PC			ENO L, EN1L, EN2L = (TS4•F SET• $\overline{\text{MS DIS}}$)•(FE + FD)	11H6	6C8	TS4	
			CPMA	CPMA LOAD = TP4•$\overline{\text{MALC}}$	10H4	6F3	TP4	
1→F			FETCH	OPR + IOT•F•CPMA LOAD = F L	12B7	12B4		

PDP-8E MECh. ZATION TABLE
FETCH FOR OPERATE GROUP TWO INSTRUCTIONS

Function	Source Req'd	Route Req'd	Dest'n Req'd	Enabling Levels & Boolean Expression	Source Page & Code	Dest'n Page & Code	Time	Major State	Instruction
MA+1→PC	CPMA			ENOL, EN1H, EN2H = TS1 (BF + KC)	11H7	6C8	TS1	FETCH	ALL
		+1		CARRY IN L = TS1 (BF + KC)	11H1	9			
			PC	PC LOAD = TP1 (BF + KC)	10J5	6B2	TP1		
0→SKIP	0			0→SKIP = BF • TP1	10H1	10J5			
SA→MD	SA			TIME STROBE = (STROBE • FIELD• WRITE) DELAYED	5B8	5E 1-8	TS2		
			MD LINES	MD DIR L = 100 NS INTO TS1, 0→MD DIR	14C7	5E 1-8			
AC→DATA	AC			AC→BUS = BF • TS2	11J3	6C6			
MD→MB	MD LINES			ENOL, EN1L, EN2H = (DCA•E•TS2)• BTS2• (MS DIS + BRK DATA CONT)	11H5	6C8			
			MB	MB LOAD = BTP2	10H2	6C2	TP2		
MD0-2→ IR			IR	LOAD IR = TP2•F	12E5	12E5			
MD 4(0) AC→DATA	AC			AC→BUS = (MD 7(1) •OPE•TS3)•(BTS3• MD4(0) •OPR•F)•AC→BUS INH	11H3	6C6	TS3		NOP
MD5(1)• AC0(1) 1→SKIP	1			SKIP DATA IN L = OP2•TS3•MD5(1) • AC0(1) •MD 8(0)	10H2	10J2			SMA
MD 6(1)• AC = (Ø) 1→SKIP	1			SKIP DATA IN L = OP2•TS3•MD 6(1) • AC = 0•MD 8(0)	10H2	10J2			SZA
MD 7(1)• L = 1 1→SKIP	1			SKIP DATA IN L = OP2•TS3•MD 7(1)• LINK (1)•MD 8(0)	10H2	10J2			SNL
MD 8(1) SKIP→ SKIP	0			SKIP DATA IN H = [OP2•TS3•MD 8(1) • (MD6(1) •AC=0)+(MD5(1) •AC 0(1) + (LINK(1) •MD 7(1))]	10H2	10J2			SPA SNA SZL SKP
	1			SKIP DATA IN L = OP2•TS3.MD 8(1) • (MD 6(1) •AC = 0 + MD 5(1) •AC0(1) + LINK(1) •MD 7(1))	10H2	10J2			
MD 9(1) SR→DATA	SR			SR→BUS = IR0•IR1•IR2•MD3(1)•MD 9(1)• MD 11(0)•USER MODE•TS3	2 D7	2B 3-5			OSR
MD 10(1) C→RUN	0			RUN DATA IN L = MD0(1)•MD1(1)•MD 2(1)• MD 3(1)•MD 10(1)•MD 11(0)USER MODE•F	14B6	14D6			HLT
DATA→AC	DATA			DATA T L, DATA F H = (OPR+IOT)•BTS3• (C2L + C1L + C0H)	11H8	6D6			ALL

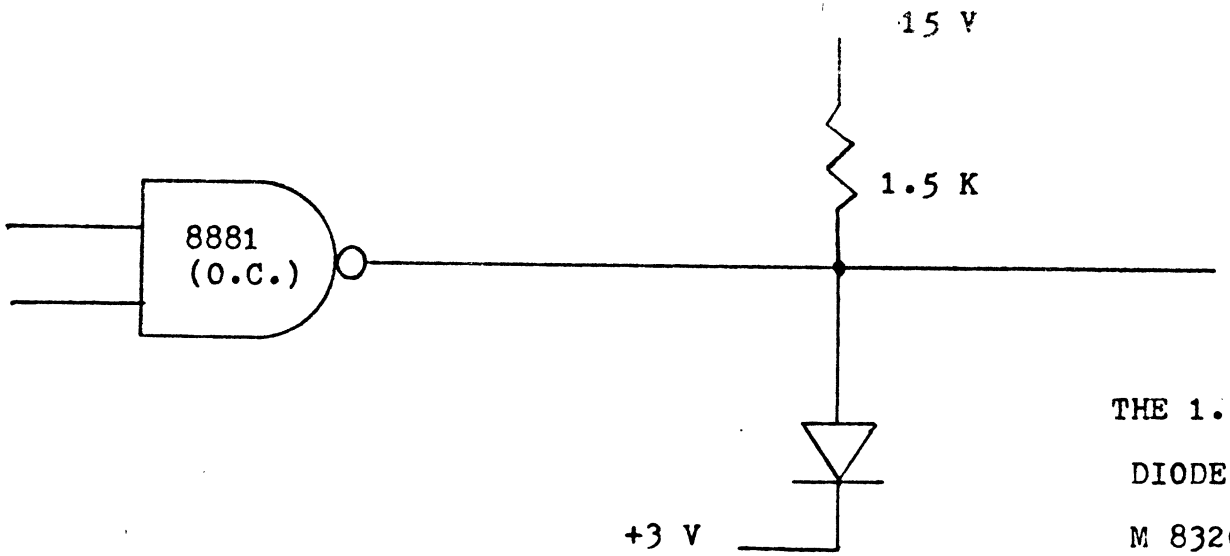
			AC	AC LOAD = OPR·F·BTP3	10H3	6B2	TP3	ALL
			SKIP	CLOCK SKIP = BTP3	10J2	10J2		ALL
SKIP PC→MA	PC			EN0L, EN1L, EN2L=(TS4·F SET·$\overline{MS\ DIS}$)·($\overline{FE+FD}$)	11H6	6C8	TS4	
SKIP PC+1→MA	PC			EN0L, EN1L, EN2L,=(TS4·F SET·$\overline{MS\ DIS}$)·($\overline{FE+FD}$)	11H6	6C8		
		+1		CARRY IN=SKIP(1)·TS4·F SET·$\overline{MS\ DIS}$)	11H1	9		
			CPMA	CPMA LOAD = TP4·$\overline{MALC}$	10H4	6F3	TP4	
1→F			FETCH	OPR+IOT·F·CPMA LOAD = FL	12B7	12B4		

PDP-8E MECHANIZATION TABLE
FETCH FOR OPERATE GROUP THREE INSTRUCTIONS EXCLUDING EAE

Function	Source Req'd	Route Req'd	Dest'n Req'd	Enabling Levels & Boolean Expression	Source Page & Code	Dest'n Page & Code	Time	Major State	Instuction
MA+1→PC	CPMA			ENOL, EN1H, EN2H = TS1(BF+KC)	11H7	6C8	TS1	FETCH	ALL
		+1		CARRY IN L = TS1(BF+KC)	11H1	9			
			PC	PC LOAD = TP1(BF+KC)	10J5	6B2	TP1		
0→SKIP	0			0→SKIP = BF·TP1	10H1	10J5			
SA→MD	SA			TIME STROBE=(STROBE·FIELD·$\overline{\text{WRITE}}$) DELAYED	5B8	5E1-8	TS2		
			MD LINES	MD DIR L = 100NS INTO TS1, 0→MD DIR	14C7	5E1-8			
AC→DATA	AC			AC→BUS L = BF·TS2	11J3	6C6			
MD→MB	MD LINES			ENOL, EN1L, EN2H = ($\overline{\text{DCA}}$·E·TS2)·BTS2·($\overline{\text{MS DIS}}$+BK DATA CONT)	11H5	6C8			
			MB	MB LOAD = BTP2	10H2	6C2	TP2		
MD0-2→IR			IR	LOAD IR = TP2·F	12E5	12E5			
MD 4(0)·MD 7(0) AC→DATA	AC			AC→BUS=(BTS3·OPR·F·MD4(0)·($\overline{\text{OPE}}$·TS3·$\overline{\text{MD7(1)}}$·$\overline{\text{AC→BUS INH}}$	11J3	6C6	TS3		NOP
MD 5(1) MQ→DATA	MQ			MQ→BUS=OPE·TS3·MD5(1)·$\overline{\text{MQ→BUS INH}}$	11J4	6C6			MQA
MD 7(1)·MD 4(0) AC→MQ	AC			SHL+Ld ENA=BTS3·OPR·F·MD 4(0)	11H3	6B5			MQL
MD 7(1)·MD 4(1) 0→MQ	0			$\overline{\text{SHL+Ld ENA}}$ = BTS3·MD4(1)·OPR·F	11H3	6B5			CLA, MQL
DATA→AC	DATA			DATA T, L, DATA F H = OPR+IOT·BTS3·(C2L+C1I+COH)	11H8	6D6			ALL
			AC	AD LOAD = OPR·F·BTP3	10H3	6B2	TP3		
			MQ	MQ LOAD = BTP3·MD 7(1)·OPE	10H4	6J2			MQL
PC→CPMA	PC			ENOL, EN1L, EN2L = (TS4·F SET·$\overline{\text{MS DIS}}$)·($\overline{\text{FE+FD}}$)	11H6	6C8	TS4		ALL
			CPMA	CPMA LOAD = TP4·$\overline{\text{MALC}}$	10H4	6F3	TP4		
1→F			FETCH	OPR+IOT·F·CPMA LOAD=FL	12B7	12B4			

PDP 8/e

BUS LOGIC CONCEPT
AND
724 POWER SUPPLY REQUIREME



THE 1.5K RESISTOR AND THE DIODE ARE LOCATED ON THE M 8320 (or the M 832) BUS LOADS CARD.

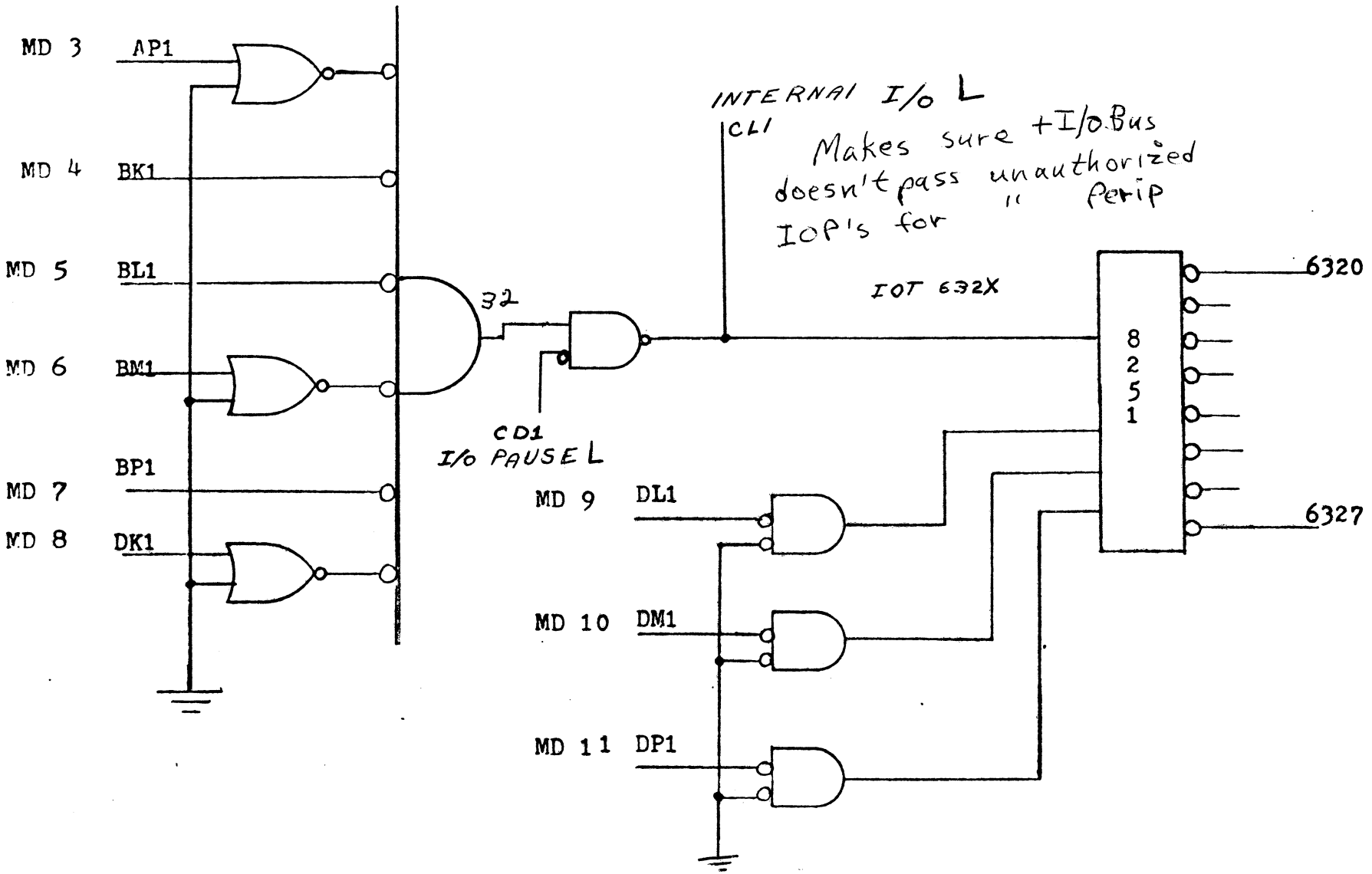
H724 POWER SUPPLY

OUTPUT VOLTAGE	+15 V	+5 V	-15 V	+8 V
CURRENT OUTPUT	1 A	20 A	.8 A	2 A
CURRENT USED	.8 A	6 A	4.5 A	1.25 A
UNUSED CURRENT	.2 A	14 A	3.5 A	

INTERR.

DEVICE CODE SELECTIO!

(632X)



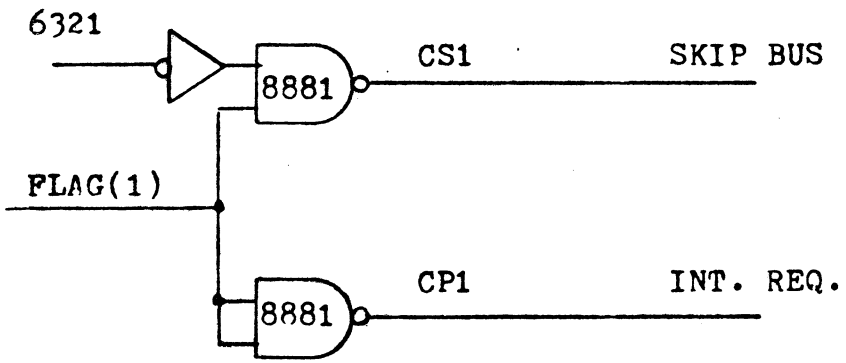
INTERNAL I/O

EXAMPLE INSTRUCTION SET

INSTRUCTION	FUNCTION
6321----------	CHECK FLAG AND SKIP IF =(1)
6322----------	SKIP ONCE IF FLAG "A" = 1, SKIP ONCE IF FLAG "B" = 1, SKIP TWICE IF FLAG "A" AND FLAG "B" ARE = 1.
6324----------	TRANSFER AC TO DEVICE
6325----------	TRANSFER AC TO DEVICE, 0---AC
6326----------	TRANSFER DATA FROM DEVICE TO AC.
6327----------	TRANSFER DATA TO THE PC.

GATING FOR SKIP AND INT. REQ.

FIG. 1

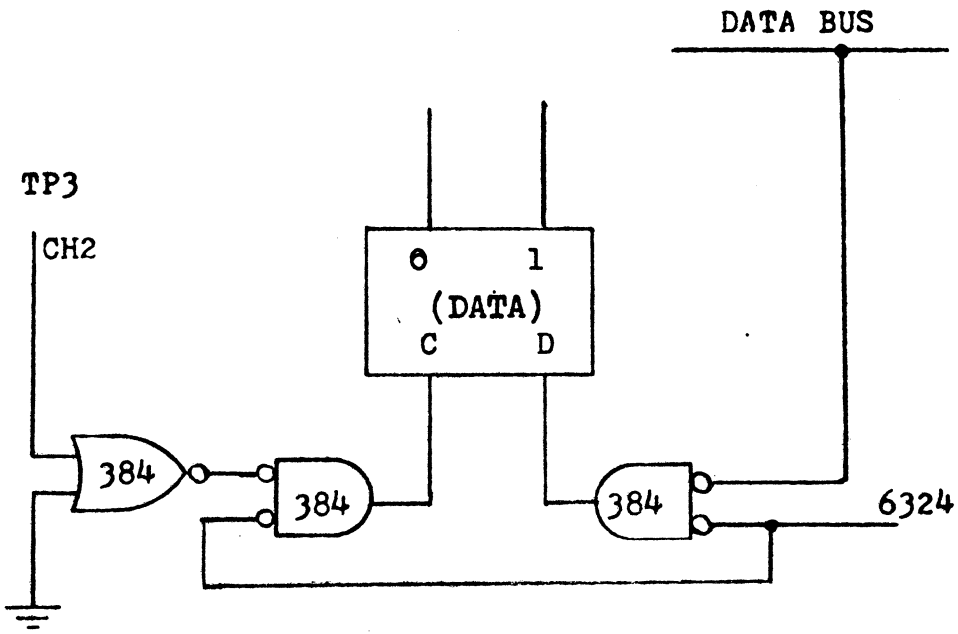


THE PURPOSE OF THIS GATING IS TO ALLOW SENSING OF SKIP CONDITIONS AND ACKNOWLEDGING INTERRUPT REQUESTS FROM THE DEVICE.

I/O OUTPUT TRANSFER GATING

C0 C1 C2
H H H = AC∨DATA----AC

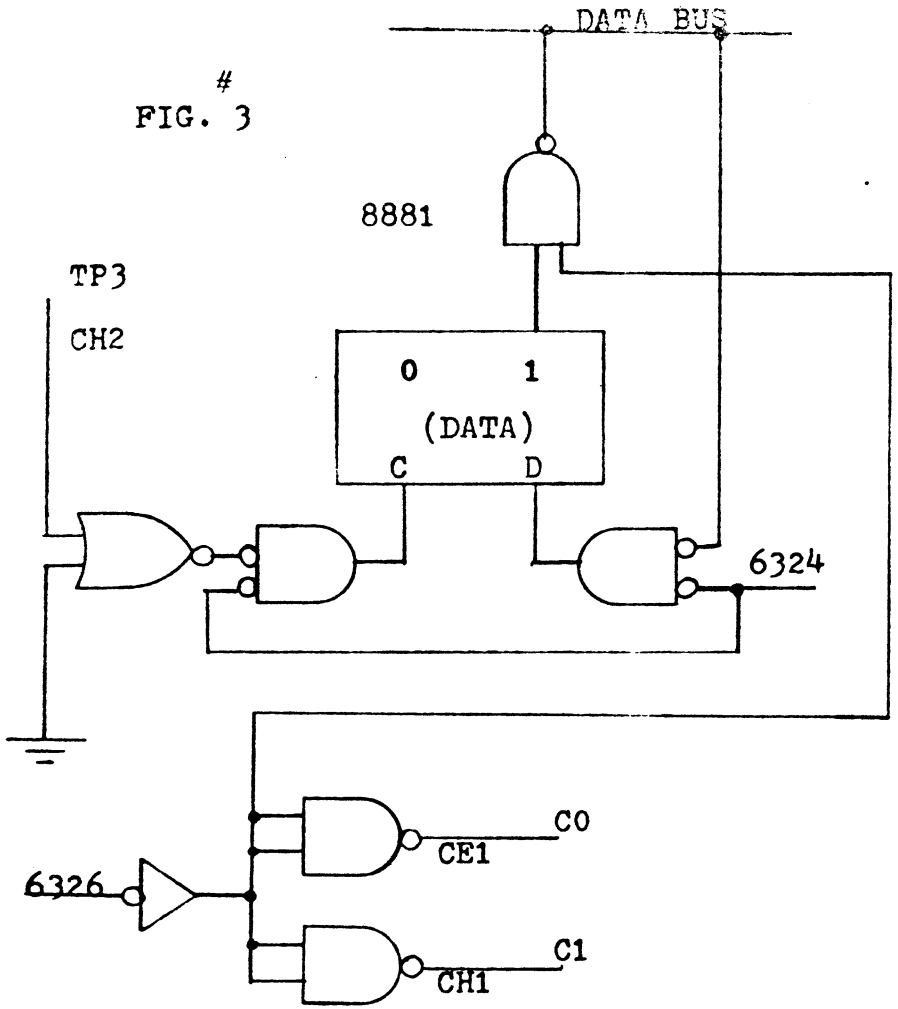
FIG. 2



THE PURPOSE OF THIS GATING IS TO ALLOW THE TRANSFER OF THE AC TO A REGISTER IN THE DEVICE.

DATA----AC GATING

FIG. #3

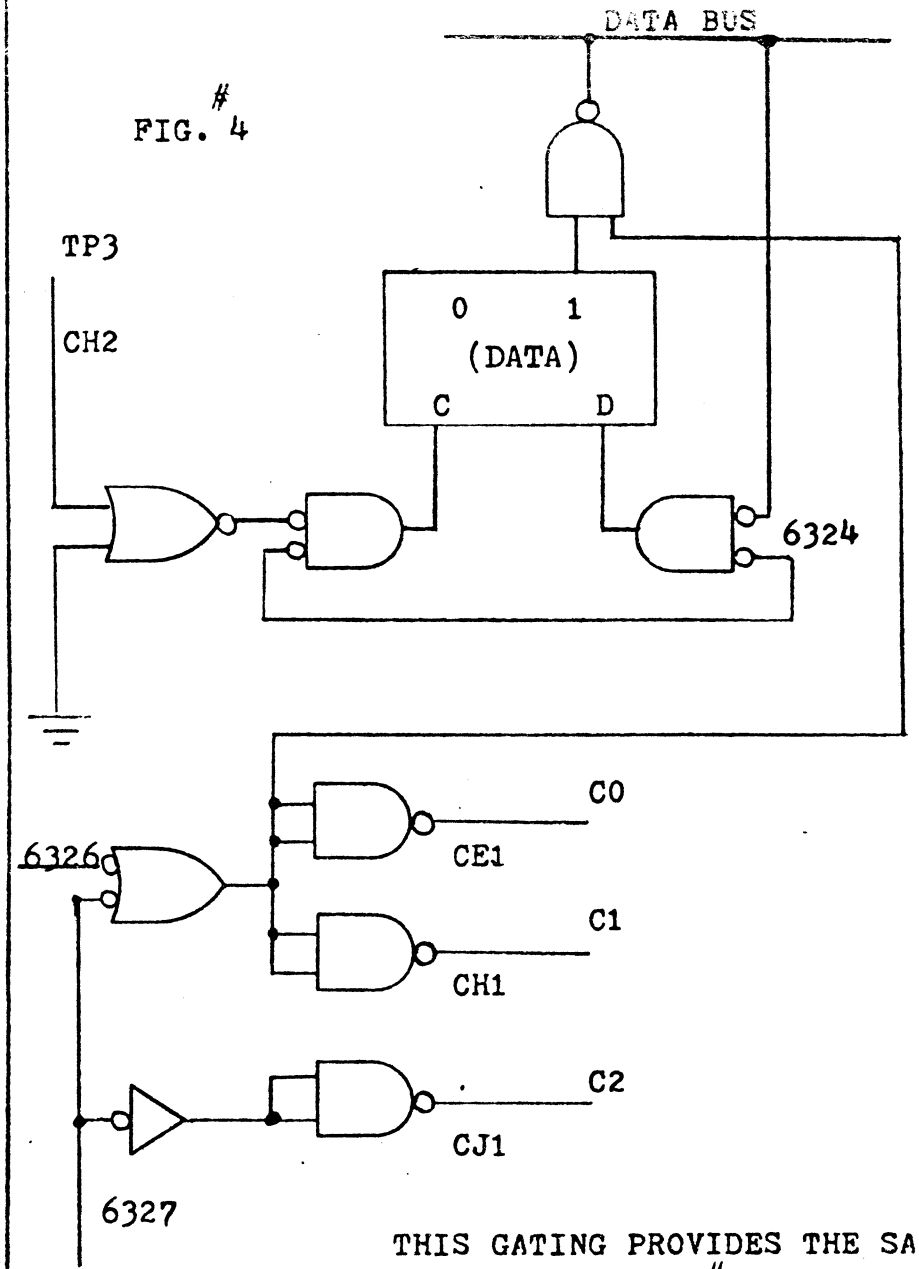


THIS GATING HAS THE SAME FUNCTION AS THAT OF FIG.#2 PLUS THE DATA FROM THE DEVICE REGISTER CAN BE TANSFERED TO THE AC.

DATA⟶AC = Co=L, C1=L, C2=H

DATA----PC GATING

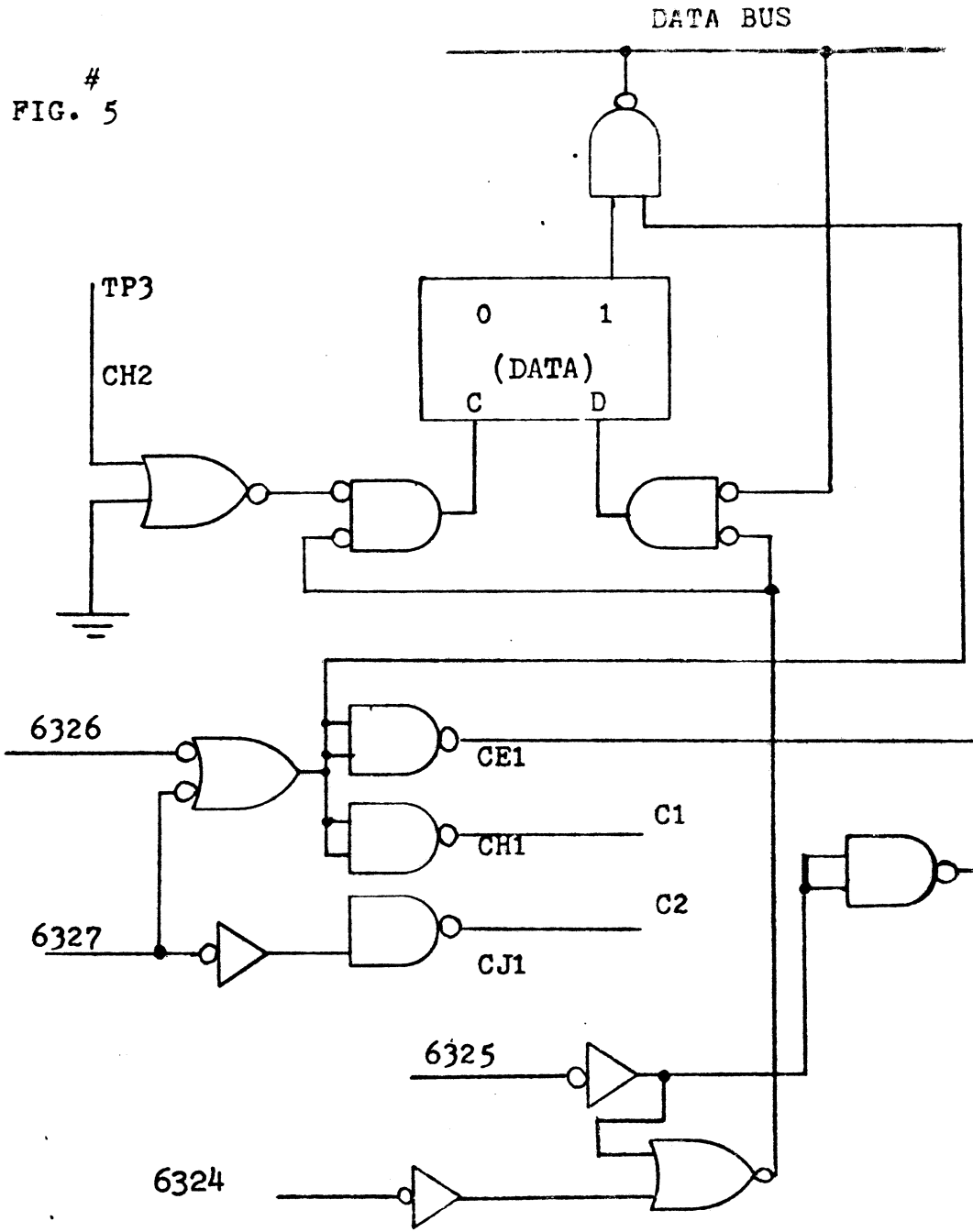
FIG. #4



THIS GATING PROVIDES THE SAME FUNCTIONS AS THOSE OF FIG. #2 and 3 PLUS THE TRANSFERING OF DATA⟶PC.

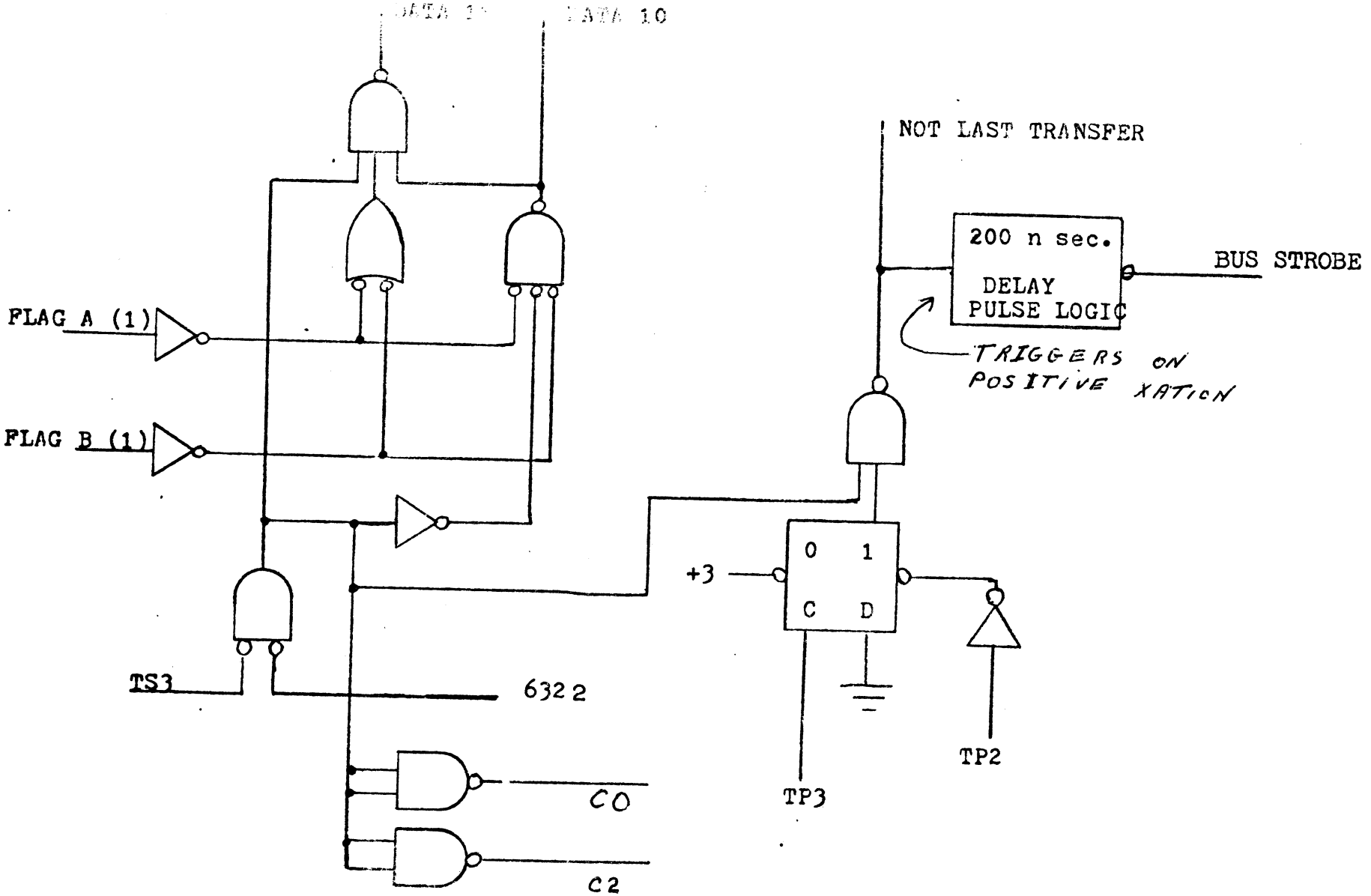
DATA⟶PC= CO=L, C1=L, C2

FIG. # 5



THIS GATING PROVIDES THE SAME FUNCTIONS AS THOSE OF FIG. # 2,3, and 4 PLUS THE 0⟶AC AFTER THE AC HAS BEEN TRANSFERED TO THE DATA BUS.(0→AC @ BUS STROBE

AC⟶DATA, 0⟶AC
CO=L, C1=H, C2=H


DATA 10
NOT LAST TRANSFER
200 n sec.
DELAY
PULSE LOGIC
BUS STROBE
TRIGGERS ON
POSITIVE XATION
FLAG A (1)
FLAG B (1)
+3
0
1
C
D
TS3
6322
TP2
TP3
CO
C2

1. BEFORE THE CPU IS ABLE TO ACKNOWLEDGE ANY INTERRUPT REQUESTS FROM DEVICES,THE INTERRUPT SYSTEM MUST BE ENABLED. THE FOLLOWING IS A PROCEDURE IN ORDER TO ENABLE THE INTERRUPT SYSTEM.

   a) AN ION INSTRUCTION,6001,MUST BE PERFORMED. THIS INSTRUCTION WILL "TURN ON" THE INTERRUPT SYSTEM ONLY HALF WAY.

   b) ANY INSTRUCTION THAT FOLLOWS THE 6001 WILL FULLY ENABLE THE INTERRUPT SYSTEM.

2. WHEN THE INTERRUPT SYSTEM HAS BEEN ENABLED, THE CPU IS CAPABLE OF ACKNOWLEDGING INTERRUPT REQUESTS. THE FOLLOWING IS A GENERAL DESCRIPTION OF ACKNOWLEDGING AN INTERRUPT REQUEST:

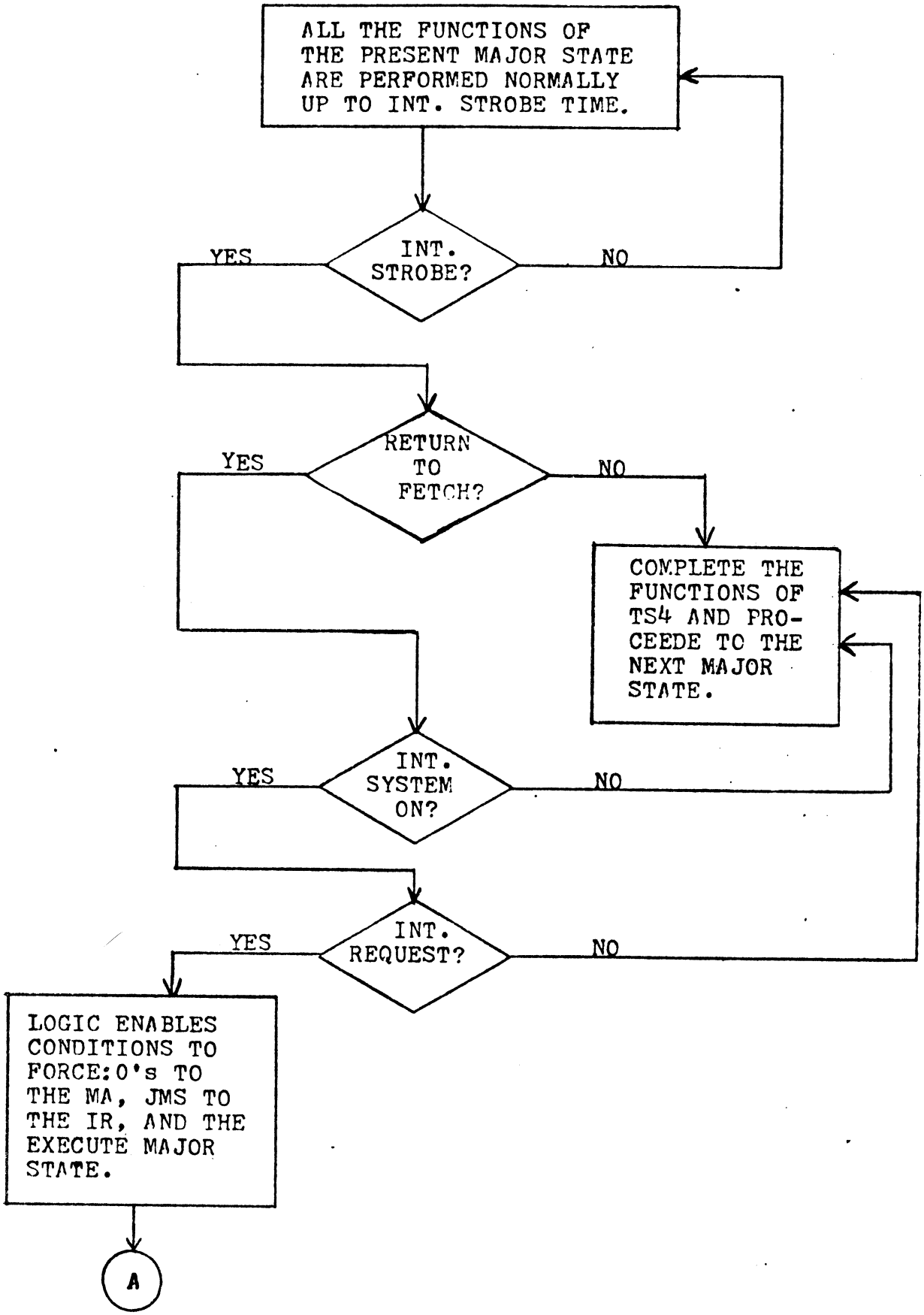
   a) THE CPU WILL ACKNOWLEDGE AN INTERRUPT REQUEST AT TP3 TIME OF ANY MAJOR STATE PROVIDING THAT FETCH IS THE NEXT MAJOR STATE TO OCCUR. BY ACKNOWLEDGING THE INTERRUPT REQUEST THE NORMAL OCCURENCES OF TS4 ARE DISABLED. INSTEAD OF PERFORMING PC---MA (OR PC + 1----MA FOR A SKIP CONDITION) 0's ARE FORCED TO THE MA. THEREFORE; THE NEXT ABSOLUTE ADDRESS TO BE REFERENCED IS 0000. ALONG WITH FORCING THE CPU TO GO TO ADDRESS 0000, THE IR IS FORCED TO DECODE A JMS INSTUCTION.

NOW THE CPU IS FORCED TO PERFORM A JMS INSTRUCTION AT ADDRESS 0000. SINCE THE CPU's IR IS FORCED TO DECODE A JMS AND ABSOLUTE ADDRESS 0000 WAS FORCED TO THE MA, THE CPU CAN ELIMINATE THE NEED FOR A FETCH CYCLE. THE CONTROLING LOGIC FOR CHANGING MAJOR STATES IS FORCED TO PRODUCE THE CONDITIONS THAT WILL ALLOW THE MAJOR STATE OF EXECUTE TO OCCUR NEXT. THE VALUE OF THE PC (WHICH IS THE NEXT INSTRUCTION'S ADDRESS FOR THE MAIN PROGRAM) WILL BE STORED IN ADDRESS 0000. ADDRESS 0000 IS NOW THE ENTRANCE POINT FOR RETURNING TO THE MAIN PROGRAM. THE FIRST INSTRUCTION OF THE SUBROUTINE WILL BE PERFORMED AT ADDRESS 0001. AT THIS POINT AN INTERROGATION ROUTINE IS EXECUTED (SEE HANDOUT #2s). AS THE JMS IS EXECUTED,THE INT. SYS. IS AUTOMATICALY SHUT OFF.

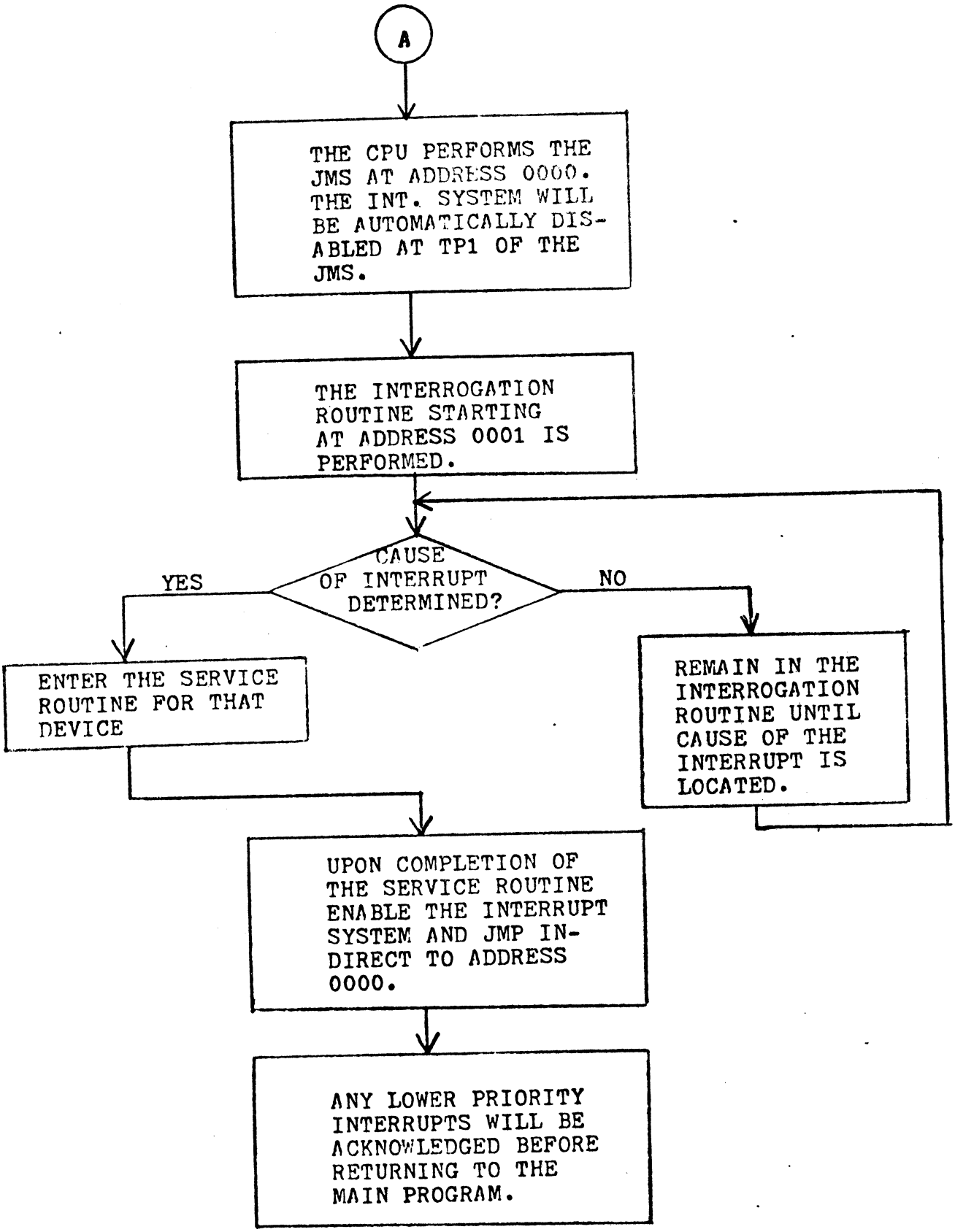
3. UPON THE COMPLETION OF THE INTERROGATION AND THE SERVICE ROUTINES, PROCEDURES FOR RETURNIG TO THE MAIN PROGRAM ARE AS FOLLOWS:

   a) BECAUSE THE INTERRUPT SYSTEM WAS DISABLED DURING THE JMS, WHICH ALLOWED THE INTERRUPT BEING PROCESSED NOT TO BE INTERRUPTED BY ANOTHER INTERRUPT, THE INTERRUPT SYSTEM SHOULD BE ENABLED TO ACKNOWLEDGE ANY LOWER PRIORITY INTERRUPTS BEFORE RETURNING TO THE MAIN PROGRAM. THIS IS ACCOMPLISHED BY PERFORMING AN ION

INSTRUCTION AT THE NEXT TO THE LAST ADDRESS OF THE SERVICE ROUTINE PROGRAM. THE ION INSTRUCTION IS THEN FOLLOWED BY THE JMP INDIRECT TO ADDRESS 0000. IF ANY LOWER PRIORITY INTERRUPTS ARE PRESENT AT THIS TIME, THEY WILL BE ACKNOWLEDGED (ACCORDING TO PRIORITY) AND THE ORIGINAL CONTENTS OF ADDRESS 0000, WHICH HOLDS THE PC OF THE FIRST INTERRUPT WILL NOT BE CHANGED.


ALL THE FUNCTIONS OF THE PRESENT MAJOR STATE ARE PERFORMED NORMALLY UP TO INT. STROBE TIME.
INT. STROBE?
YES
NO
RETURN TO FETCH?
YES
NO
COMPLETE THE FUNCTIONS OF TS4 AND PROCEEDE TO THE NEXT MAJOR STATE.
INT. SYSTEM ON?
YES
NO
INT. REQUEST?
YES
NO
LOGIC ENABLES CONDITIONS TO FORCE:0's TO THE MA, JMS TO THE IR, AND THE EXECUTE MAJOR STATE.
A

INTERRUPT SEQUENCE FLOW DIAGRAM



INTERROGATION ROUTINE

EXAMPLE PROGRAM

```
0000/ PC (XXXX)
0001/ DCA 3050
0002/ GTF 6004
0003/ DCA 3051
0004/ JMP I 5405
0005/ 7000

7000/ 6XXX
7001/ SKP 7410
7002/ JMP 5250 — DEVICE
7003/ 6XXX
7004/ SKP 7410
7005/ JMP 5300 — DEV.
7006/ 6XXX
7007/ SKP 7410
7010/ JMP 5350 — DEV.
```

DESCRIPTION OF BASIC INTERROGATION ROUTINE

THE ABOVE PROGRAM IS AN EXAMPLE OF A BASIC INTERROGATION ROUTINE. IN ADRESS 0000 THE PC (WHICH IS THE ADDRESS OF THE NEXT INSTRUCTION FOR THE MAIN PROGRAM) IS STORED FROM THE FORCED JMS OF THE INTERRUPT SYSTEM. THE DCA 3050 LOCATED IN ADDRESS 0001 IS USED TO STORE THE AC VALUE IN CASE THE DEVICE CAUSING THE INTERRUPT TRANSFERS INFORMATION TO THE AC. THE CONTENTS OF ADDRESS 0002, A GTF, IS USED TO BRING THE VALUE OF THE LINK INTO THE AC. THE GTF IS FOLLOWED BY A DCA 3051 , IN LOCATION 0003, SO THAT THE VALUE OF THE LINK MAY BE STORED IN MEMORY. THE JMP I LOCATED IN ADDRESS 0004 IS USED TO EXIT PAGE "0". THE REASON FOR EXITING PAGE"0" IS THAT IT IS USED FOR AUTO INDEX PURPOSES AS WELL AS DIRECT ACCESS FROM ALL PAGES.

THE JMP I TAKES THE PROGRAM TO ADDRESS 7000 WHERE THE INTERROGATION ROUTINE BEGINS. THE FIRST PART OF THE PROGRAM WAS WRITTEN FOR "HOUSE CLEANING" PURPOSES. THE INSTRUCTION LOCATED IN ADDRESS 7000, LISTED AS 6XXX, WILL CHECK INTERRUPT FLAG OF THE FIRST DEVICE (PRIORITY 0) FOR A SKIP CONDITION. IF THAT DEVICE'S INTERRUPT FLAG WAS SET, A SKIP CONDITION WOULD RESULT CAUSING THE PROGRAM TO REFERENCE ADDRESS 7002 INSTEAD OF 7001. THE INSTRUCTION AT ADDRESS 7002, AJMP 5250, WILL ENTER THE SERVICE ROUTINE FOR DEVICE PRIORITY 0.

EXAMPLE PROGRAM

IF DEVICE 0'S INTERRUPT FLAG WAS NOT SET, THE PROGRAM WOULD REFERENCE ADDRESS 7001. AN UNCONDITIONAL SKIP, 7410, LOCATED AT ADDRESS 7001 WILL CAUSE THE PROGRAM TO SKIP OVER THE JMP 5250. THUS THE ENTRANCE POINT FOR THE SERVICE ROUTINE OF DEVICE 0 IS SKIPPED OVER.

THE SAME SEQUENCE OF CHECKING THE INTERRUPT FLAGS APPLIES TO THE OTHER TWO DEVICES.

THE NEXT PORTION OF THE EXAMPLE PROGRAM DEMONSTRATES HOW TO ENTER THE MAIN PROGRAM AND ALSO ENABLE THE INTERRUPT SYSTEM. THIS PORTION OF THE PROGRAM WILL BE EXECUTED AFTER THE MAIN PORTION OF THE SERVICE ROUTINE HAS BEEN COMPLETED.

(AS AN EXAMPLE, DEVICE PRIORITY 1 WILL BE USED)

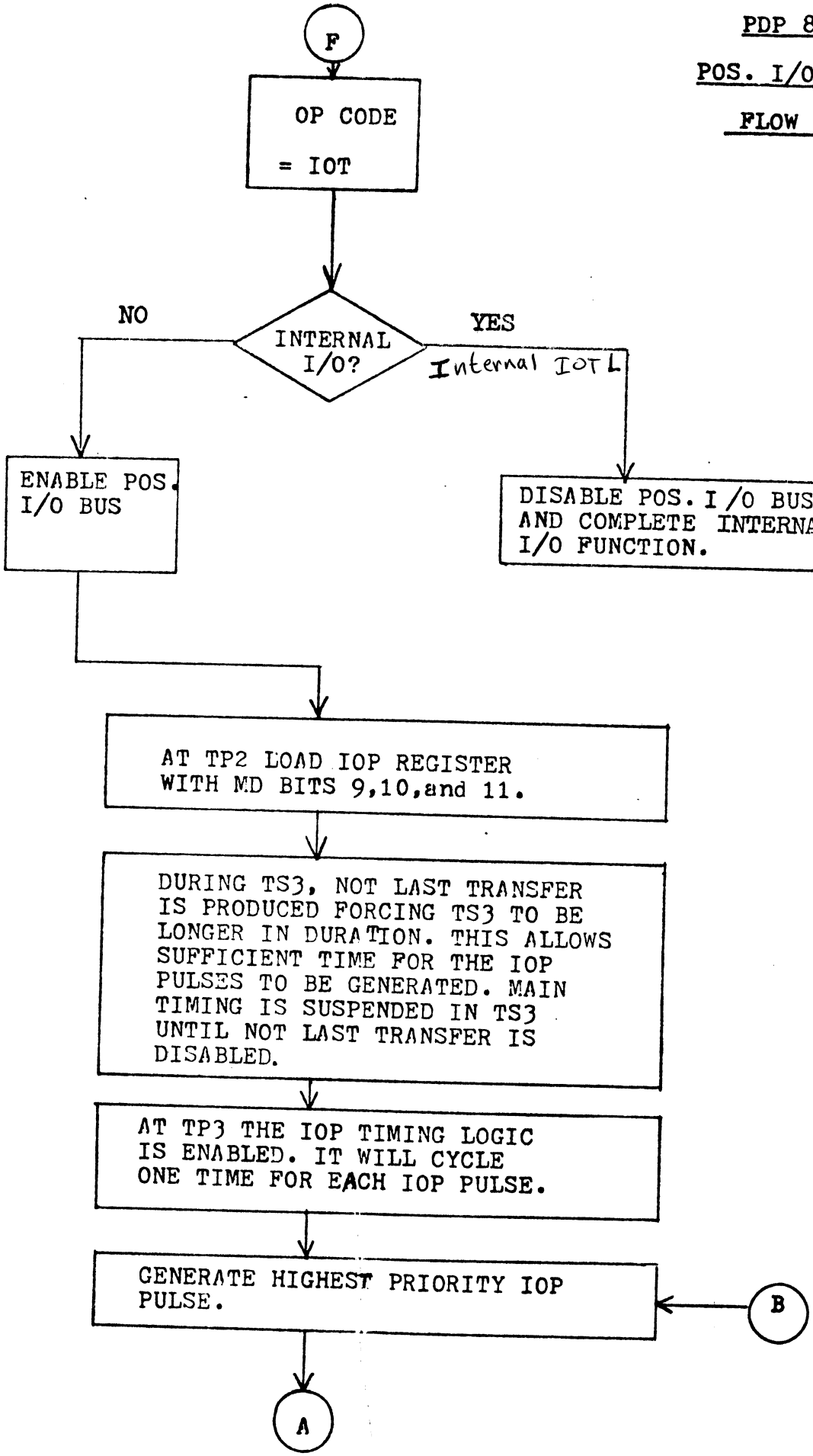
```
7124/ CLA  7200
7125/ TAD  1051
7126/ RTF  6005
7127/ CLA  7200
7130/ TAD  1050
7131/ ION  6001
7132/ JMP I 5400
```

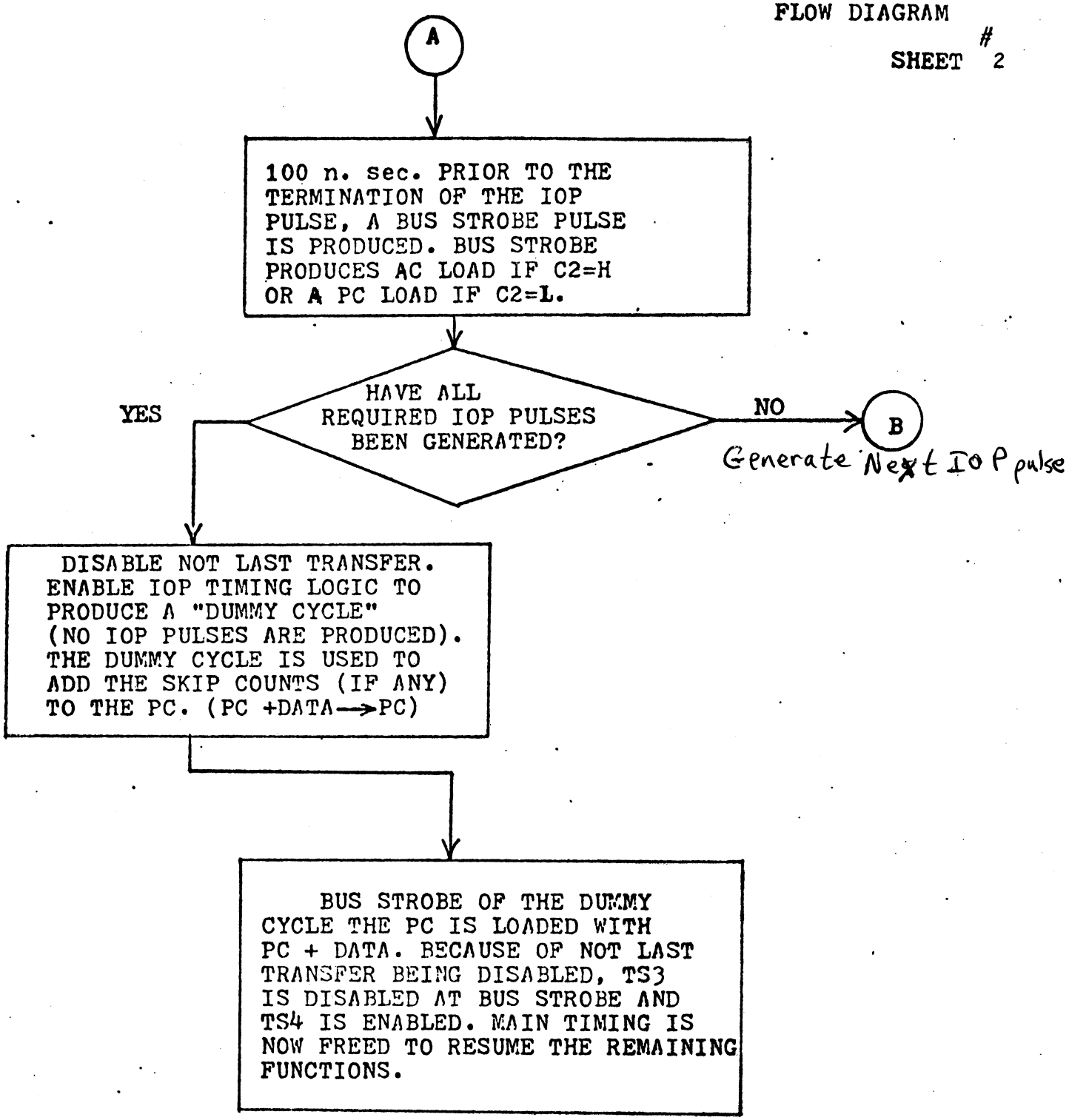
AT LOCATION 7124, A CLA CLEARS THE AC IN PREPARATION FOR RESTORING THE VALUES THAT WERE STORED IN MEMORY BY THE "HOUSE CLEANING" PORTION OF THE PROGRAM. THE TAD INSTRUCTION, 1051, LOCATED AT ADDRESS 7125 WILL BRING THE ORIGINAL CONDITION OF THE LINK, PRIOR TO THE INTERRUPT, INTO THE AC. FOLLOWING THE TAD INSTRUCTION IS A RTF INSTRUCTION. THIS INSTRUCTION WILL RETURN THE ORIGINAL CONTENTS OF THE LINK, WHICH IS NOW IN THE AC, BACK INTO THE LINK. THE NEXT LOCATION, 7127, HOLDS A CLA INSRUCTION. THIS WILL BE USED TO CLEAR THE AC IN PREPARATION FOR RESTORING THE ORIGINAL AC, PRIOR TO THE INTERRUPT. AFTER THE AC HAS BEEN CLEARED BY THE CLA, THE TAD INSTRUCTION, 1050, AT ADDRESS 7130 WILL BRING THE ORIGINAL CONTENTS OF THE AC, PRIOR TO THE INTERRUPT, BACK INTO THE AC. THE 6001 INSTRUCTION, ION, WILL PARTIALLY ENABLE THE INTERRUPT SYSTEM. THE LAST INSTRUCTION, 5400, WILL FULLY ENABLE THE INTERRUPT SYSTEM AND ALLOW THE PROGRAM TO RETURN TO ADDRESS 0000. FROM THIS POINT THE PROGRAM CAN ENTER THE MAIN PROGRAM OR ACKNOWLEDGE ANY LOWER PRIORITY INTERRUPTS THAT MAY BE PRESENT.

PDP 8/e

POS. I/O BUS

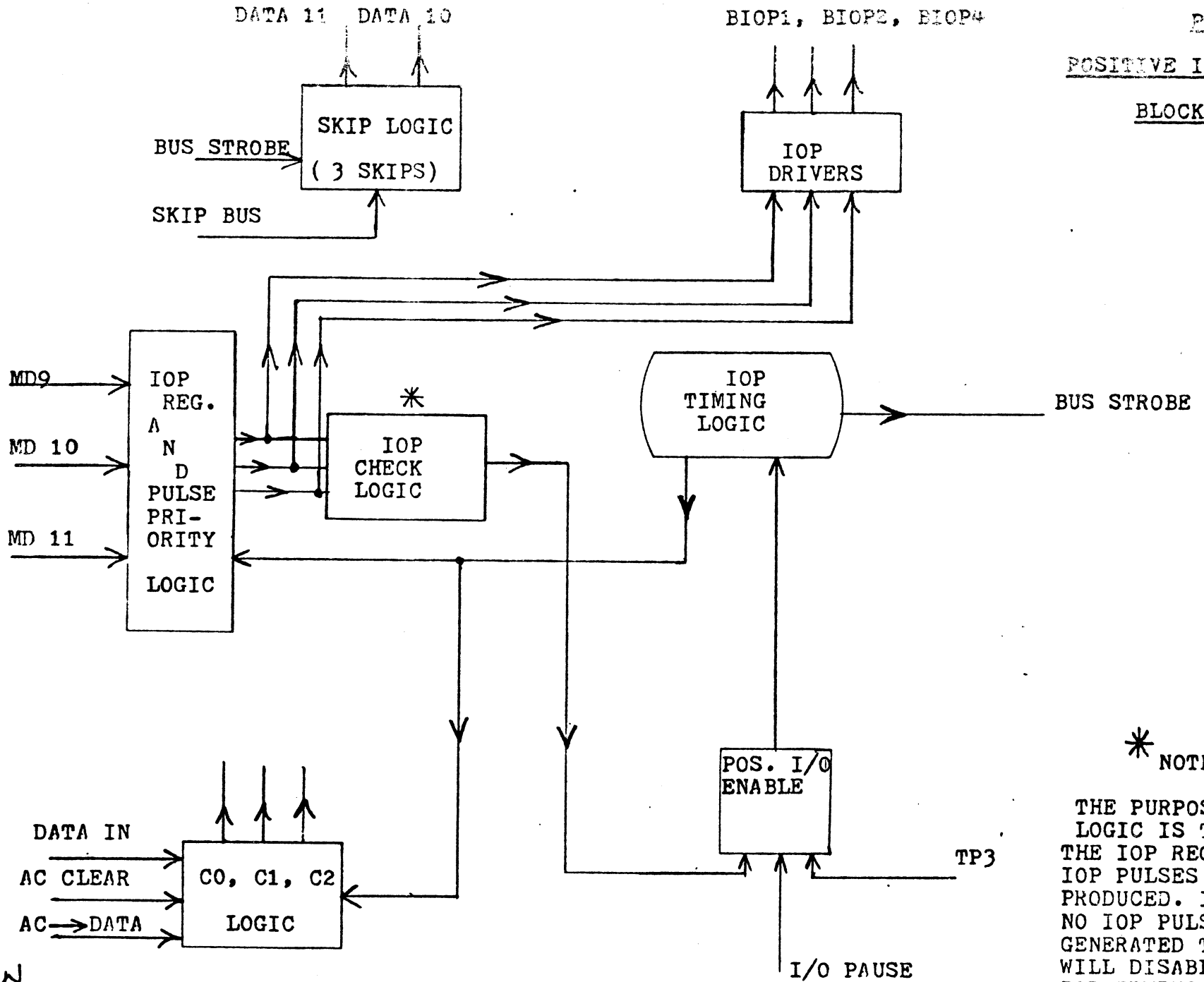
FLOW DIAGRAM




A
100 n. sec. PRIOR TO THE TERMINATION OF THE IOP PULSE, A BUS STROBE PULSE IS PRODUCED. BUS STROBE PRODUCES AC LOAD IF C2=H OR A PC LOAD IF C2=L.
HAVE ALL REQUIRED IOP PULSES BEEN GENERATED?
YES
NO
B
Generate Next IOP pulse
DISABLE NOT LAST TRANSFER. ENABLE IOP TIMING LOGIC TO PRODUCE A "DUMMY CYCLE" (NO IOP PULSES ARE PRODUCED). THE DUMMY CYCLE IS USED TO ADD THE SKIP COUNTS (IF ANY) TO THE PC. (PC +DATA→PC)
BUS STROBE OF THE DUMMY CYCLE THE PC IS LOADED WITH PC + DATA. BECAUSE OF NOT LAST TRANSFER BEING DISABLED, TS3 IS DISABLED AT BUS STROBE AND TS4 IS ENABLED. MAIN TIMING IS NOW FREED TO RESUME THE REMAINING FUNCTIONS.
BIT 11
IOP 1
600 NS TO 3μS
200 NS TO 1μS
BIT 10
IOP 2
BIT 9
IOP 4
100 NS
Should be low going signal
BUS STROBE

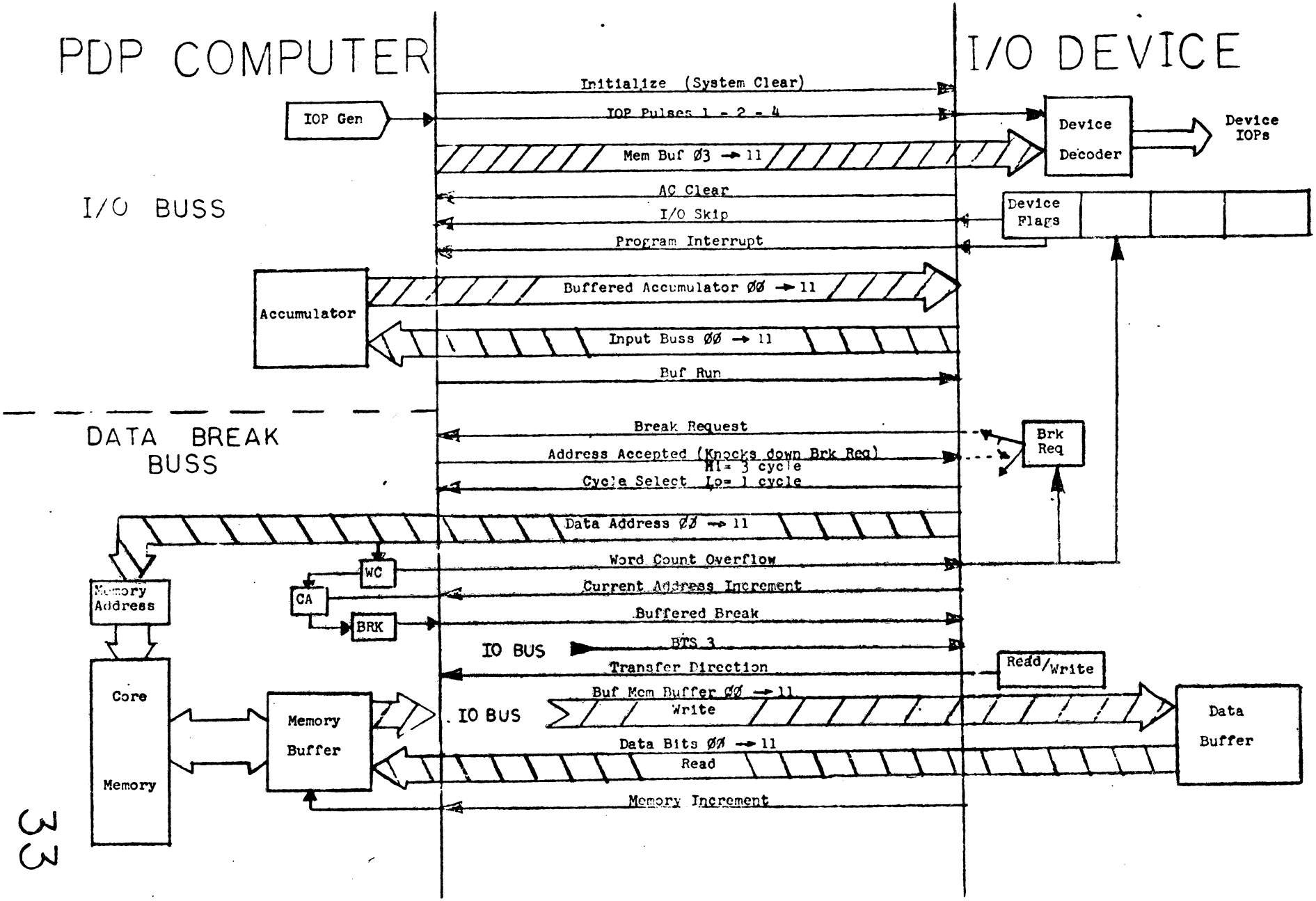
POSITIVE I/O INTERFACE

BLOCK DIAGRAM



* NOTE:

THE PURPOSE OF THIS LOGIC IS TO CHECK THE IOP REGISTR FOR IOP PULSES TO BE PRODUCED. IF THERE ARE NO IOP PULSES TO BE GENERATED THIS LOGIC WILL DISABLE THE IOP TIMING LOGIC..


PDP COMPUTER
I/O DEVICE
Initialize (System Clear)
IOP Gen
IOP Pulses 1 - 2 - 4
Device Decoder
Device IOPs
Mem Buf Ø3 → 11
I/O BUSS
AC Clear
I/O Skip
Device Flags
Program Interrupt
Accumulator
Buffered Accumulator ØØ → 11
Input Buss ØØ → 11
Buf Run
DATA BREAK BUSS
Break Request
Brk Req
Address Accepted (Knocks down Brk Req)
H1 = 3 cycle
Cycle Select Lo = 1 cycle
Data Address ØØ → 11
WC
Word Count Overflow
Memory Address
CA
Current Address Increment
BRK
Buffered Break
IO BUS
BTS 3
Transfer Direction
Read/Write
Core Memory
Memory Buffer
Buf Mem Buffer ØØ → 11
IO BUS
Write
Data Buffer
Data Bits ØØ → 11
Read
Memory Increment

STATE FUNCTIONS OF 3 CYCLE BREAK

WORD COUNT

THE PURPOSE OF THIS STATE IS TO INCREMENT A MEMORY WORD LOCATED AT A KNOWN MEMORY ADDRESS. THIS KNOWN MEMORY ADDRESS IS SPECIFIED BY THE 3 CYCLE DEVICE. THE MEMORY WORD LOCATED AT THIS ADDRESS IS THE 2'S COMPLIMENTED VALUE OF THE NUMBER TRANSFERS TO BE PERFORMED.

IF THE MEMORY WORD IS INCREMENTED TO 0000, WORD COUNT OVERFLOW IS SENT TO THE DEVICE TO INHIBIT ANY MORE BREAK REQ.

CURRENT ADDRESS

THE PURPOSE OF THIS STATE IS TO SPECIFY AN ADDRESS AT WHICH THE TRANSFER TAKES PLACE.

THIS FUNCTION CAN BE DONE 2 WAYS. 1) A KNOWN VALUE, LOCATED AT THE INCREMENTED VALUE OF THE MEMORY ADDRESS OF THE WORD COUNT STATE WILL BE INCREMENTED EACH TIME THIS STATE IS REFERENCED. THIS INCREMENTED VALUE IS USED FOR THE MA OF THE BREAK STATE, THUS ALLOWING SEQUENTIAL DATA TRANSFERS FROM MEMORY. 2) THE KNOWN VALUE NEED NOT BE INCREMENTED ALLOWING DATA TRANSFERS AT A CONSTANT ADDRESS. HOWEVER; THIS METHOD REQUIRES A BACKGROUND PROGRAM.

BREAK

THE PURPOSE OF THIS STATE IS TO PERFORM THE TRANSFERS. THERE ARE 4 FUNCTIONS WHICH MAY BE ACCOMPLISHED 1) DATA MAY BE TRANSFERED FROM THE DEVICE TO THE CPU. 2) DATA FROM THE CPU MAY BE TRANSFERED TO THE DEVICE. 3) DATA FROM THE DEVICE MAY BE ADDED WITH A MEMORY WORD FROM THE CPU. 4) A MEMORY WORD OF THE CPU MAY BE INCREMENTED.

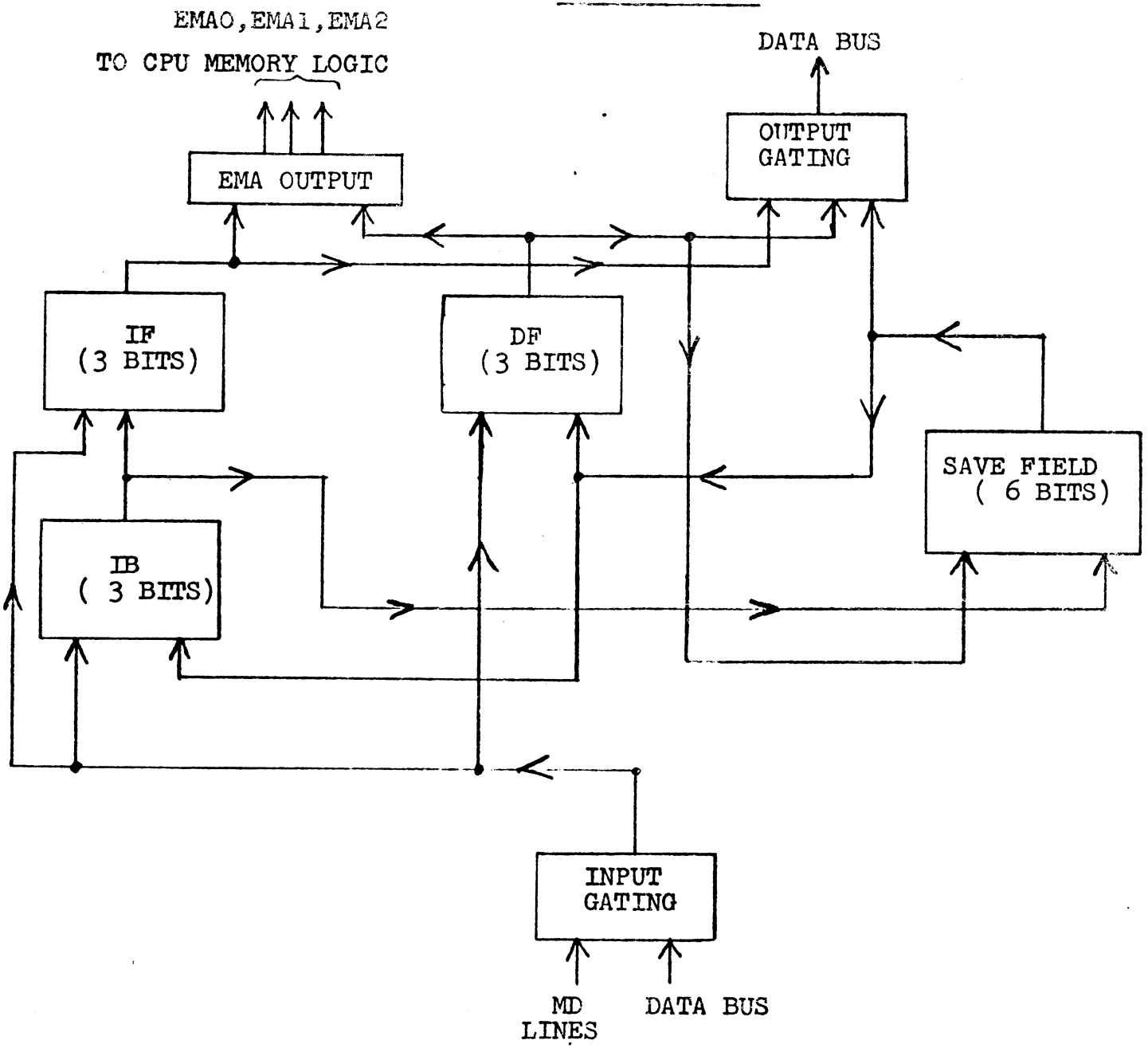
ALL 4 FUNCTIONS ARE PERFORMED AT A MEMORY ADDRESS SPECIFIED BY THE CURRENT ADDRESS STATE.

OPTIONS FOR BREAK MAJOR STATE

WRITE (DATA OUT) DI (0) INC (0)	(MEM) ⟶ MDL ⟶ PERIPH.	CONTENTS OF MEMORY SENT TO PERIPH AND WRITTEN BACK INTO MEMORY
READ (DATA IN) DI (1) INC (0)	DATA LINES ⟶ DATA BUS ⟶ MB MDL ⟶ (MEM)	DATA FROM PERIPH WRITTEN INTO MEMORY
INC (MEM) DI (0) INC (1)	(MEM) ⟶ MDL DATA BUS = 1 DATA BUS + MDL ⟶ MB ⟶ (MEM)	CONTENTS OF MEMORY +1 WRITTEN BACK INTO MEMORY (USED AS A COUNTER)
DATA ADDED TO DI (1) INC (1)	(MEM) ⟶ MDL DATA LINES ⟶ DB DATA BUS + MDL ⟶ MB ⟶ (MEM)	BREAK DATA IS ADDED TO THE CONTENTS OF MEMORY AND WRITTEN BACK INTO MEMORY

EXTENDED MEMORY CONTROL

BLOCK DIAGRAM



FUNCTIONS OF INSTUCTION FIELD AND DATA FIELD REGISTERS

THE FOLLOWING INFORMATION DESCRIBES THE RELATIONSHIP BETWEEN THE CPU MAJOR STATES AND EXTENDED MEMORY FUNCTIONS.

KEY FUNCTIONS: ALTHOUGH THE CPU IS IN NO MAJOR STATE, EXTENDED MEMORY MUST STILL CONTROL WHAT FIELD WILL BE REFERENCED DURING KEY FUNCTIONS. BECAUSE OF THIS, THE PROGRAMMER HAS A CHOICE OF ANY FIELD (LIMITED, OF COURSE, TO THE AMOUNT OF CORE ON THE SYSTEM) TO MANUALLY DEPOSIT OR EXAMINE DATA. THE CONTENTS OF THE IF REGISTER DETERMINES WHICH FIELD WILL BE REFERENCED DURING KEY FUNCTIONS.

FETCH: THE CONTENTS OF THE IF REGISTER WILL DETERMINE WHAT FIELD WILL BE REFERENCED TO OBTAIN INSTRUCTIONS. THIS CONDITION IS TRUE FOR ANY FETCH MAJOR STATE.

DEFER: THE CONTENTS OF THE IF REGISTER WILL DETERMINE FROM WHAT FIELD THE POINTER ADDRESS WILL BE OBTAINED.* HOWEVER; AT TP4 TIME OF DEFER EITHER THE IF OR THE DF REGISTER CAN BE THE REGISTER THAT WILL SPECIFY THE FIELD

AT TP4 TIME OF EACH MAJOR STATE EXT. MEM. CONT. LOGIC DECIDES IF THE IF OR DF WILL SPECIFY THE NEXT FIELD TO REFERENCE.

INSTRUCTION FIELD AND DATA FIELD

REGISTERS

DEFER---------- (cont.) TO BE REFERENCED FOR THE EXECUTE MAJOR STATE. THIS IS DEPENDENT UPON WHAT TYPE OF INSTRUCTION IS BEING PERFORMED.

IF AN AND, TAD, ISZ, OR DCA IS BEING PERFORMED, THE CONTENTS OF THE DF REGISTER WILL DETERMINE WHICH FIELD WILL BE REFERENCED TO OBTAIN DATA FOR THE EXECUTE MAJOR STATE. IF THE CONTENTS OF THE DF ≠ IF THE PROGRAMMER NOT ONLY HAS THE CAPABILITY OF OBTAINING DATA FROM ANY MEMORY PAGE; HE ALSO HAS THE CAPABILITY OF REFERENCING A DIFFERENT MEMORY FIELD IN ORDER TO OBTAIN DATA. ON THE OTHER HAND; IF THE DF = IF, A NORMAL INDIRECT IS PERFORMED IN THE SAME FIELD THE INSTRUCTION AND THE POINTER WERE OBTAINED.

IF A JMS OR JMP IS BEING PERFORMED, THE DF REGISTER WILL NOT SPECIFY THE FIELD TO BE REFERENCED IN EXECUTE; THE IF REGISTER WILL.THE REASON FOR THIS IS THAT A JMP NEVER ENTERS THE EXECUTE STATE. IF THE DF REGISTER COULD SPECIFY WHICH FIELD TO REFERENCE DURING THE NEXT MAJOR

FUNCTIONS OF

INSTRUCTION FIELD AND DATA FIELD

REGISTERS

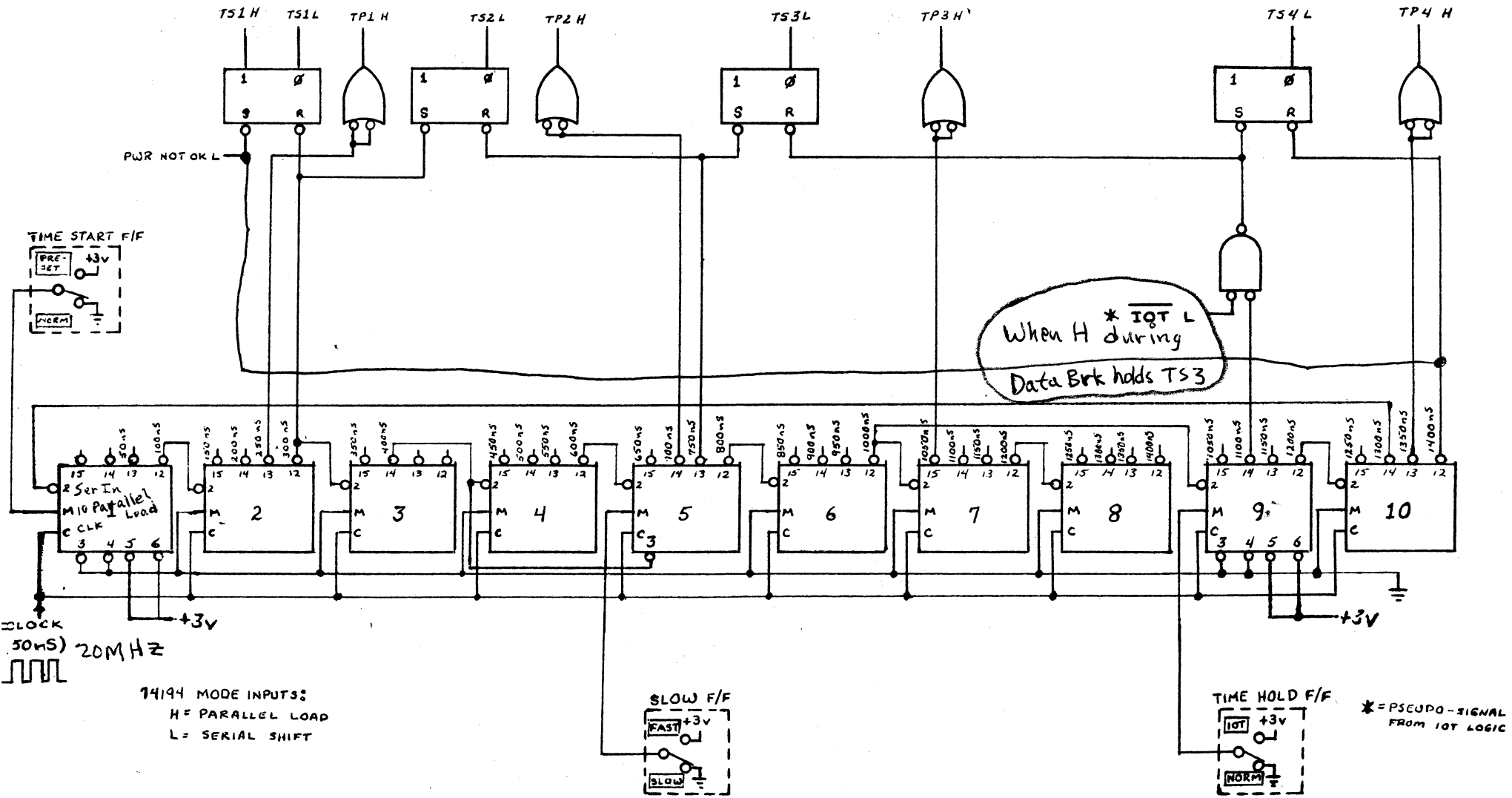
DEFER ---------- (cont.)

STATE, WHICH WOULD BE FETCH, THE NEXT INSTRUCTION WOULD BE OBTAINED FROM THE FIELD SPECIFIED BY THE DF REGISTER. THIS WOULD DEFEAT THE PURPOSE OF THE DF REGISTER AS IT IS USED ONLY TO OBTAIN DATA; NOT INSTRUCTIONS.

BECAUSE THE JMP INDIRECT CANNOT USE THE DF REGISTER TO SPECIFY A FIELD TO REFERENCE FOR THE EXECUTE MAJOR STATE, THE JMS INSTRUCTION IS FORCED TO DO THE SAME. IF THE JMS INSTRUCTION COULD USE THE DF REGISTER TO SPECIFY WHICH FIELD TO REFERENCE FOR THE EXECUTE MAJOR STATE, THE PC WOULD BE STORED IN THAT PARTICULAR FIELD. HOWEVER; UPON RETURNING FROM THE SUBROUTINE, WHICH REQUIRES THE USE OF A JMP INDIRECT TO THE ADDRESS WHERE THE PC IS LOCATED, THE FIELD WHERE THE PC IS STORED COULD NOT BE REFERENCED. WHY? A JMP INDIRECT CANNOT REFERENCE A FIELD SPECIFIED BY THE DF REGISTER.

EXECUTE ---------- THE EXECUTE MAJOR STATE CAN REDERENCE A FIELD SPECIFIED BY THE IF OR DF. (SEE ABOVE INFORMATION OF DEFER FOR EXPLANATION)

All Times Delayed 50ns, wait for first Clk, Pls. on 1st Time thru.



M8330 TIMING GENERATOR- SIMPLIFIED SCHEMATIC DIAGRAM