

## PUBLICATIONS REVISION

System 80

OS/3  
Dump Analysis  
Programming Guide

UP-9980 Rev. 1

This Library Memo announces the release and availability of the *System 80 OS/3 Dump Analysis Programming Guide*, UP-9980 Rev. 1.

This guide is a standard library item (SLI). It is part of the standard library provided automatically with the purchase of the product.

This guide describes how to interpret Operating System/3 (OS/3) dumps to help you understand job or system conditions. It contains the following:

- Descriptions of the structure and functions of OS/3 software
- Explanations of the task control block (TCB) and the program status word (PSW)
- Descriptions of three OS/3 dumps
- Sample dump analyses

For Release 12.0, this guide adds parameters to the suboption SELECT and incorporates minor technical and editorial changes.

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**UNISYS**

**System 80  
OS/3**

**Dump Analysis  
Programming Guide**

OS/3 Release 12.0

October 1988

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Printed in U S America  
UP-9980 Rev. 1



**UNISYS**

**System 80  
OS/3**

**Dump Analysis  
Programming Guide**

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# About This Guide

## Purpose

This guide discusses dump analysis primarily as it applies to user jobs.

## Scope

This guide contains the following:

- Explanations of the task control block (TCB) and the program status word (PSW), two of the most important structures in dump analysis, and an overview of the structure and functions of Operating System/3 (OS/3) software
- An overview of OS/3 dumps
- User dump analysis with examples

## Audience

The intended audience for this guide is the programmer, who will use dump analysis to help develop and maintain Unisys software, programs, and files.

## Prerequisites

Anyone using this guide should understand the following:

- Unisys OS/3 machine language
- How to read a dump
- Job organization

## How to Use This Guide

Read the entire guide to familiarize yourself with the basic concepts it presents; then use it for reference as needed.

## Organization

This document contains four sections and two appendixes:

### **Section 1. Introduction to Dump Analysis**

Defines dump analysis, and explains when you need dumps.

### **Section 2. OS/3 Overview**

Gives a basic overview of OS/3, including the supervisor and user region, and several data structures and registers.

### **Section 3. Dumps and Their Formats**

Describes the three types of dumps OS/3 provides: \$YSDUMP, JOBDUMP, and EOJ dump.

### **Section 4. Sample Dump Analyses**

Presents examples of dump analyses applied to actual user programs by using the ideas presented in the previous sections.

### **Appendix A. Program Exceptions**

Lists the interrupt codes and their causes.

### **Appendix B. \$YSDUMP File Allocation**

Shows the number of cylinders required for \$YSDUMP file allocation, depending on your system's main storage capacity and the type of disk device you are using.

## Results

After reading this document, programmers will be able to use dump analysis to correct program errors.

## Related Product Information

The following Unisys documents may be useful in understanding and implementing dump analysis.

**Note:** *Throughout this guide, when we refer you to another manual, use the version that applies to the software level at your site.*

***Processor - System 80 Programmer Reference (UP-8881)***

Provides the detailed hardware-oriented information required to program the System 80 processor. Includes a brief functional description and configurations.

***Supervisor Technical Overview (UP-8831)***

Provides an overview of the OS/3 supervisor and its functions for OS/3 high-level language programmers and site administrators.

***Assembler Programming Guide (UP-8913)***

Describes, for both novice and experienced programmers, the OS/3 basic assembly language (BAL) and its use. Discusses general language concepts, assembler instructions, and programming techniques.

***System Messages Reference Manual (UP-8076)***

Lists the system messages and describes them. Message description composed of the remedial action or response required as applicable.

***Job Control Programming Guide (UP-9986)***

Provides information on the format and use of job control statements and job control procedure calls (jprocs).

***Hardware and Software Programming Quick-Reference Guide (UP-8868)***

Summarizes machine instructions, supervisor related information, physical input/output control system (PIOCS) information, and I/O sense data byte definitions.

***System Service Programs (SSP) Operating Guide (UP-8841)***

Describes the system services programs, the utility programs that support the operation and organization of the operating system.

***Operations Guide (UP-8859)***

Describes the hardware configuration of each System 80 model, presents procedures for initializing the system, and describes all commands and procedures used within the OS/3 environment.

***Installation Guide (UP-8839)***

Provides the system administrator with information and procedures necessary to install, tailor, and maintain OS/3 software in a System 80 environment.

## Notation Conventions

The notation conventions used in this guide are the following:

- Capital letters, commas, equal signs, and parentheses must be coded exactly as shown. For example:

```
SET SY,LOFF
```

```
(SELECT)
```

```
FILE=$YSDUMP
```

- Lowercase letters and words are generic terms representing information that must be supplied by the user. Such lowercase terms may contain hyphens and acronyms (for readability). For example:

vsn (volume serial number)

decimal-area-number

job/symbiont-name

did (device identification number)

- Information contained within braces represents mandatory entries of which one must be chosen, such as:

```
DISPLAY= { startaddr:n  
          X'xxxxxxxx' [startaddr-endaddr]  
          C'ccccccc'  
          CR  
          IO  
          PR  
          RR  
          SN  
          SR  
          JOBS }
```

- Information contained within brackets represents optional entries that (depending upon program requirements) are included or omitted. Braces within brackets signify that one of the specified entries must be chosen if that parameter is to be included. For example:

```
[,P=did]
```

```
[,V= { vsn  
      { (vsn,AT) } } ]
```

- An optional parameter having a list of optional entries may have a default specification that is supplied by the operating system when the parameter is not specified by the user. Although you can specify the default, you do not have to. For easy reference, when a default specification occurs in the format delineation it is printed on a shaded background. For example:

```
[ ,SPL { PRINTER }
      { TAPE } ]
```

```
[ ,FILE= { $YSDUMP }
          { lblname } ]
```

If, by parameter omission, the operating system performs some complex processing other than parameter insertion, it is explained in the parameter description.

- When a portion of a parameter is underlined, only that portion need be specified. For example:

```
[ ,MLIB=NO]
```

can be coded as:

```
[ ,MLIB=N]
```



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# Section 1

## Introduction to Dump Analysis

### 1.1. General

Dump analysis is one of the set of diagnostic tools available for your use in developing and maintaining Unisys Operating System/3 (OS/3) software, programs, and files. A raw dump itself is a hexadecimal image of the system main storage and registers at the time of the dump, which is output to a printer. Dump facilities in OS/3 supplement a raw dump by extracting from it information needed to interpret job or system conditions. Dump analysis is the process by which you interpret the dump and is one of your primary tools for uncovering errors and problems in your OS/3 programs.

### 1.2. When Dumps Are Needed

Dumps are used on a number of occasions:

- When a HALT AND PROCEED (HPR) occurs. Many HPRs listed in the *System Messages Reference Manual* (UP-8076) request that the operator perform a dump.
- When a system hang occurs. If the system remains in a run state but cannot communicate with the operator, a dump is often taken.
- When a user job terminates with an error. If certain job control language (JCL) options are specified, the system will react to the error by generating a dump and sending it to a printer.
- When some system errors will automatically generate a system dump.

While the dump is an important diagnostic tool to you, it is an equally important tool for your Unisys representative, if a system problem is serious. For less serious problems, though, you can analyze your own dumps and resolve your program problems. In fact, the emphasis in this guide is on analyzing job failures, since system failures usually require the services of a Unisys representative.



# Section 2

## OS/3 Overview

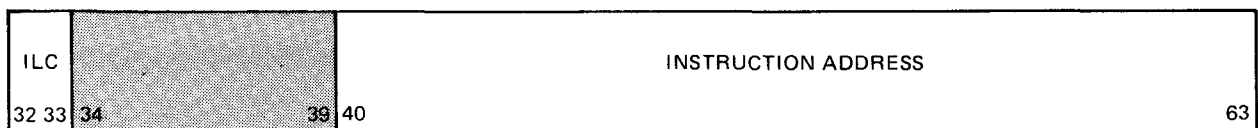
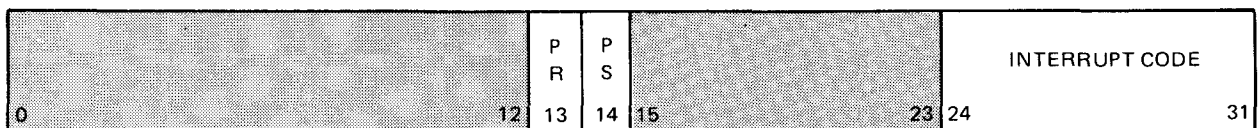
### 2.1. General

To do dump analysis, you should know some of the operation of OS/3 software. For OS/3 operation, main storage can be divided into two parts: the supervisor and the user region. The supervisor is the complex of routines that prepares user programs for running, serves as the primary I/O software interface, and, most importantly for dump analysis, handles unexpected or error conditions. The user region contains your program and is where most user-solvable problems occur. The two areas, supervisor and user region, occupy separate parts of main storage and share in the use of several data structures and registers, the most important of which are the program status word (PSW) and the task control block (TCB).

#### 2.1.1. Program Status Word (PSW)

One of the most important registers in OS/3 is the current program status word. This register contains the address of the next instruction to be executed as well as other indicators and flags which, taken together, define the hardware status of the system.

The format and description of the PSW are:



NOTE:

= Not used

A complete description of the PSW can be found in the appropriate processor programmer reference. For our discussion, you need know only about the unshaded fields shown in the previous illustration:

- PR (bit 13)

Defines the general register set to be used in executing an instruction. The processor contains two sets of 16 general registers: problem registers and supervisor registers. The set the processor uses is determined by the following settings of bit 13:

PR=1    Problem registers

PR=0    Supervisor registers (not for user analysis)

- PS (bit 14)

Specifies one of two processor modes of operation:

PS=1    Problem mode

PS=0    Supervisor mode (not for user analysis)

In problem mode, all privileged machine instructions are invalid, and their attempted execution results in a program exception (interrupt code 02; see Table A-1). In supervisor mode, all instructions may be executed.

- Interrupt code (bits 24-31)

Contains the interrupt code used in software analysis of interrupts. A list of the interrupt codes used in OS/3 is shown in Table A-1.

- ILC (bits 32-33)

Contains the instruction length code, the length in half words of the instruction being executed at the time of an interrupt. For each instruction, the instruction address field is incremented by the number of bytes indicated in the ILC for that instruction.

- Instruction address (bits 40-63)

Contains the address of the next instruction to be executed, obtained when the OS/3 hardware adds the address of the currently executing instruction to the length of that instruction.



Because the last field of the PSW, the instruction address, points to the next instruction, it can be altered by hardware branch instructions. The address can also be stored in main storage, if the current program is interrupted, and later loaded back into the PSW, permitting the program to resume without loss of control. The OS/3 software handles a number of such interrupt types, including:

- Program exception - issued by the hardware when it detects the improper use of a machine instruction or data.
- Machine check - issued by the hardware when it detects a hardware failure.
- Supervisor call - issued by user programs through use of the SVC instruction. SVC is the user's primary interface to the supervisor and may be used to call transients (2.2.1).

The OS/3 software contains error routines that handle each type of interrupt. Figure 2-1 shows how an interrupt is handled. When an interrupt occurs, the processor microcode swaps the current PSW of the interrupted program (PROG1) with a new PSW used for handling the interrupt.

The PSW for PROG1 (old PSW) is stored in a reserved slot and is used later to determine at which instruction PROG1 is to resume. The new (now current) PSW of PROG2 is loaded into the PSW register and sets the system in a new state that branches to the supervisor error routine (PROG2) that is designed to handle the interrupt.

When PROG2 routine is finished, a LOAD PSW instruction usually loads the old PSW back into the PSW register, thus returning processor control to PROG1. The hardware maintains a pair of PSWs, old and new, for each type of interrupt.

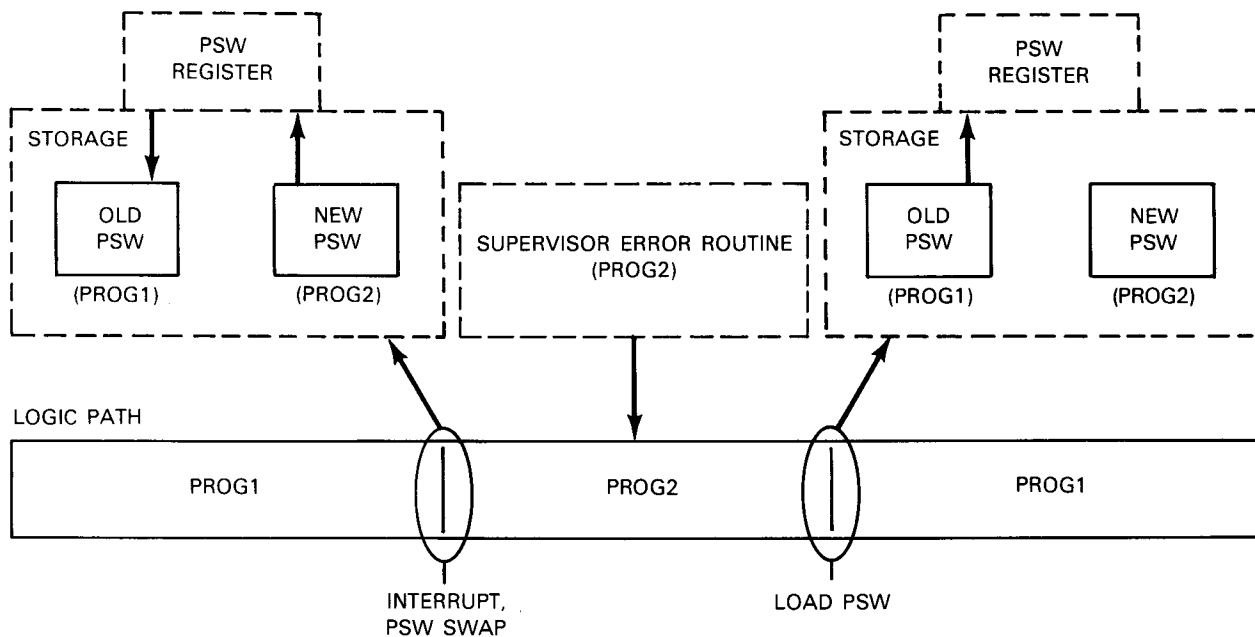


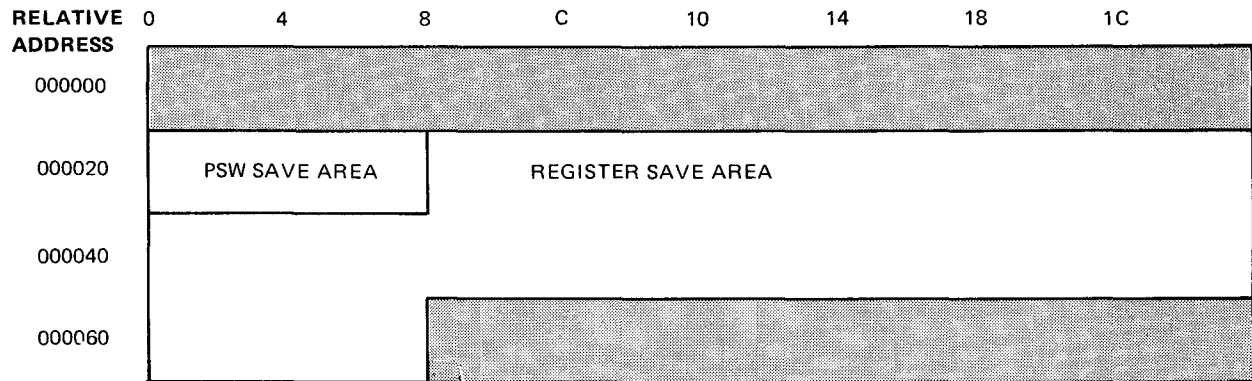
Figure 2-1. Handling an Interrupt

Because these interrupts may represent either an error condition or the operating system's response to one, their old PSWs can provide clues to the cause of an error.

### 2.1.2. Task Control Block (TCB)

Besides PSWs, the structure you may analyze most thoroughly in a dump is the task control block (TCB). Up to 255 TCBs may be active in OS/3 at any time, and each of these represents a task that is competing for processor time with all other tasks. A TCB contains fields that completely define its task to the rest of the system; the TCB thus serves as a software counterpart to the hardware-oriented PSW.

For dump analysis, certain fields are important, and these are represented by the unshaded portion of the following TCB format diagram:



NOTE:

= Not used

- PSW Save Area (byte offset  $20_{16}$ - $27_{16}$ )

Two full words into which the old PSW is stored in the event of an interrupt occurring in the task. When the task resumes processor control, this data is loaded back into the current PSW.

On abnormal termination, the PSW is also stored in the last five bytes of the preamble and the interrupt code and instruction address are displayed and written to the job log. This information is useful if a dump is not obtained.

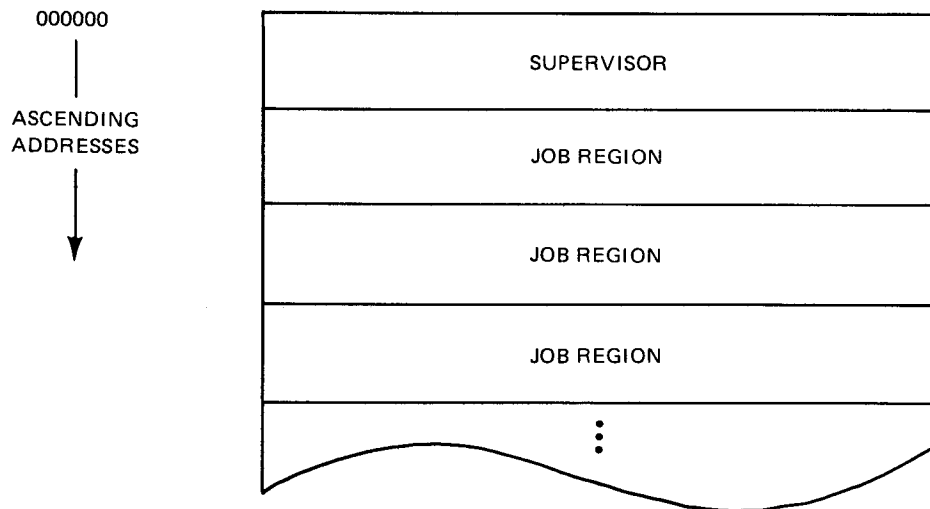
- Register Save Area (byte offset  $28_{16}$ - $67_{16}$ )

Sixteen full words into which the contents of the current register set are stored when an interrupt occurs in the task. When the task resumes processor control, these full words are loaded back into the proper register set.

You can find more detailed information of TCBs in the *Supervisor Technical Overview* (UP-8831).

## 2.2. OS/3 Structural Overview

This subsection summarizes some of the important OS/3 structures that are resident in main storage and describes how they are linked together. A knowledge of these is important to dump analysis. Broadly, main storage is allocated as follows:



The OS/3 supervisor occupies the low-order part of main storage. Job regions occupy some parts of high-order main storage, with free regions taking up the remainder. The formats of the supervisor and of the job regions are discussed in the following subsections. Samples of these are in the dumps shown in Section 3.

### 2.2.1. Supervisor

Three areas of the supervisor are often used in dump analysis: low-order main storage, the system information block, and the transient area.

## Low-Order Main Storage and Relocation Registers

The first 4096 bytes of the main memory unit for the System 80, models 8 through 20 are reserved for special purposes by the operating system and thus cannot be used as a storage area for ordinary data and instructions. In System 80, models 3 through 6, low-order main storage occupies the lowest 256 bytes of main storage. One area of low-order main storage is of interest in dump analysis; this is shown in the following illustration:

Byte Address (Hexadecimal)	0 1 2 3	4 5 6 7	8 9 A B	C D E F
00X				
01X				
02X	I/O old PSW		I/O new PSW	
03X	Exigent machine check old PSW		Exigent machine check new PSW	
04X	Program check old PSW		Program check new PSW	
05X	Supervisor call old PSW		Supervisor call new PSW	
06X	External old PSW		External new PSW	
07X	Repressible machine check* old PSW		Repressible machine check* new PSW	
08X	PER old PSW		PER new PSW	
09X	Restart old PSW*		Restart new PSW*	
0AX				

\* Not used in Model 8.

Low-order locations 000020 through 00009F contain, for each type of interrupt supported in System 80, the new PSW that is always loaded into the current PSW register, and the old PSW that is swapped out. The PSWs for program checks, machine checks, and SVCs have already been discussed; of these, the program exception will figure the most in the dump analyses shown in Section 4.

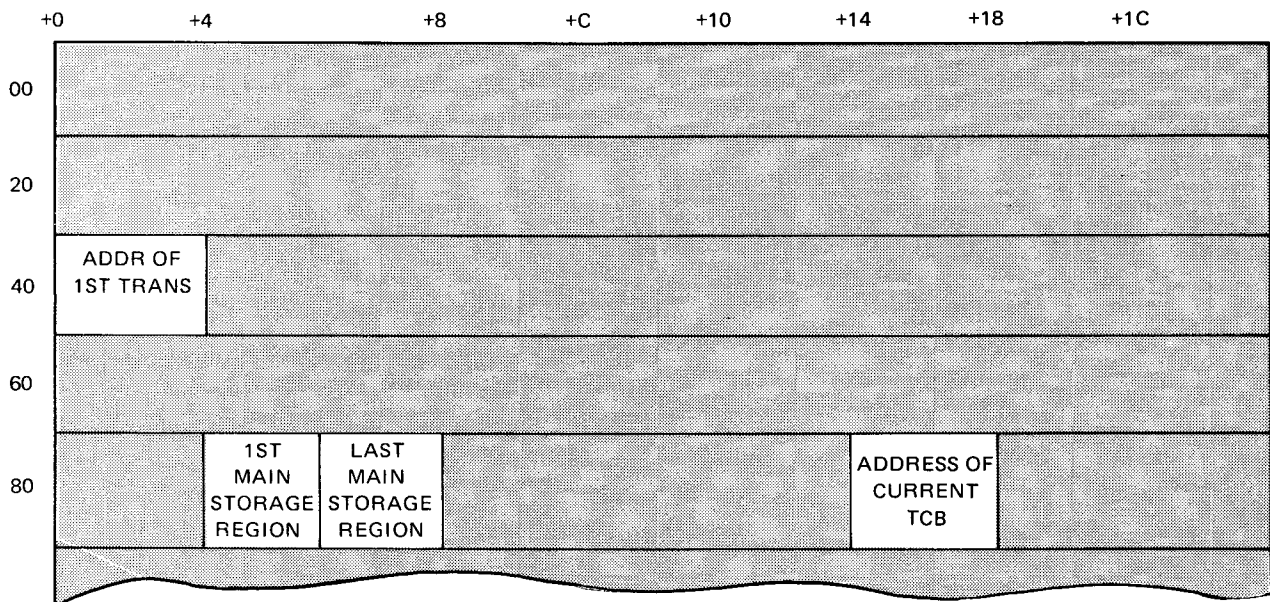
In addition to the data kept in low-order main storage, the system uses 16 relocation registers (identified as 0 through  $F_{16}$ ) for models 3 through 6 and 64 relocation registers (identified as 1- $40_{16}$ ) for models 8 through 20. In the operating system, models 8 through 20 provide up to 48 user jobs, each using one of the registers 1- $30_{16}$  as its base. For models 3 through 6, the maximum number of user jobs is 14 and each job base is one of the registers 1- $E_{16}$ .

These registers should not be confused with the base registers used within programs, especially by BAL programmers; in fact, every effective address within a job is added to the job's relocation register to get the absolute address required, a process generally transparent to users.

All TCBs within a job specify the relocation register for that job in byte 0. While the job has processor control, its register number occupies bits 8-11 of the current PSW. Register 0 is reserved for the supervisor; the field contains a value of zero, which permits supervisor routines to address the entire address space of a system. Because it is important to distinguish between system errors and job errors in dump analysis, job regions and relocation registers are further discussed in 2.2.2.

### System Information Block (SIB)

The system information block (SIB) occupies supervisor space just past low-order main storage. Fields within the SIB are often used in dump analysis and are presented in the following illustration:



NOTE: [shaded box] = Not used

Bytes 40 to 43 contain the full-word address of the first transient area (see 2.2.1). The two half words in bytes 84 to 87 define the user region with the keys to the first and last user blocks. The address in bytes 94 to 97 points to the current TCB, which is the last TCB given control by the system (see 2.3.1).

## Transients

To handle system tasks OS/3 uses two basic types of routines: supervisor overlays and transient routines. Supervisor overlays occupy a single 1192-byte area within the supervisor; these are routines that are loaded when needed to handle urgent tasks, such as operator communications. Transient routines are loaded when needed, occupy several contiguous 1192-byte areas of supervisor, and can be called by user jobs or by other transients.

### 2.2.2. Job Regions

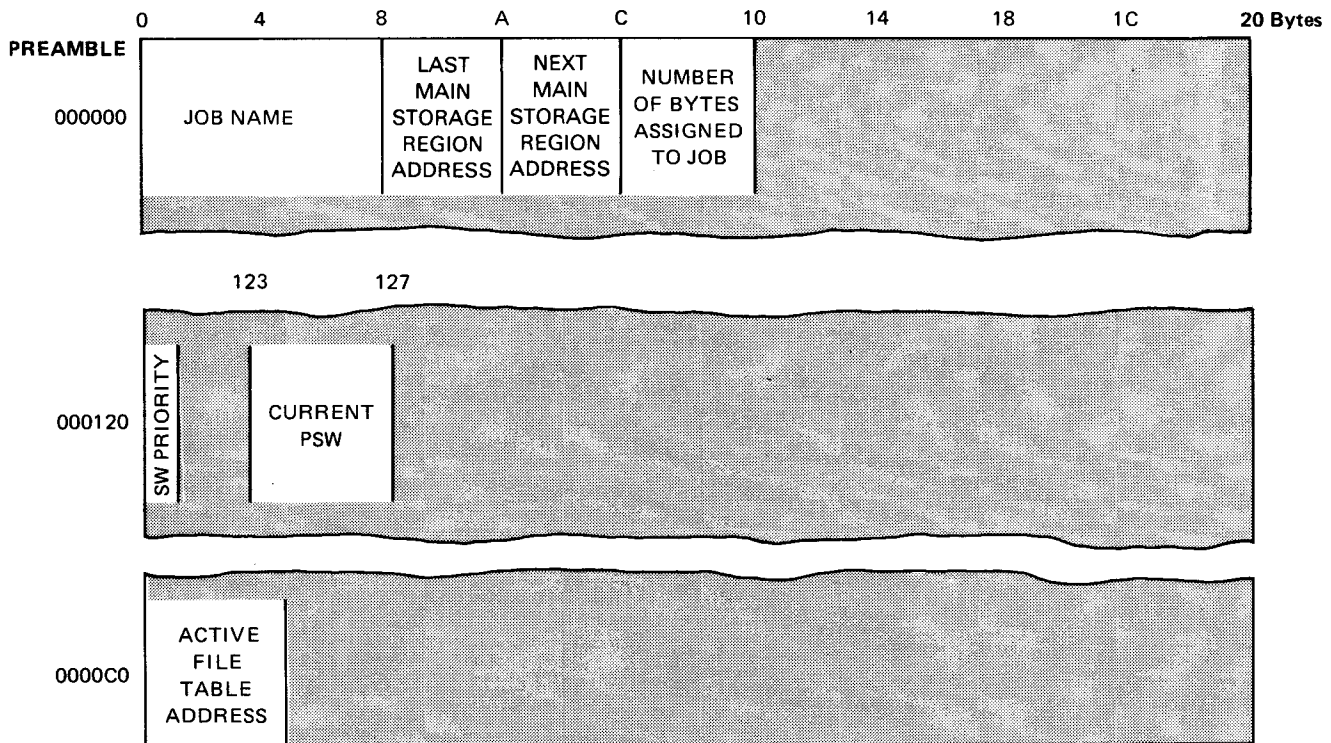
Each user job in OS/3 occupies a single job region in main storage. Within that region are all the fields unique to that job, including:

- Job preamble
- Job TCBs
- Load module


The job preamble contains data about the job that is used by the supervisor in scheduling, running, accounting, and terminating the job. The TCBs, of which at least one exists for each job, represent the individual tasks that may compete for processor time while the job is run. The load module is scheduled and loaded through job control language (JCL) or as a symbiont and may contain one or more phases.

## OS/3 Overview

As mentioned in 2.2.1., models 3 through 6 have 14 user slots and models 8 through 20 have 48 user slots. This means that System 80 allows you to run simultaneously up to 14 user jobs on models 3 through 6 and up to 48 user jobs on models 8 through 20. Pointers within each job preamble point to other jobs and to free regions (see 2.2.4). Each job uses one of the relocation keys discussed in 2.2.1. Five fields in the preamble are often used in dump analysis, as shown in the following illustration:



NOTE:

 = Not used in analysis

The eight bytes starting at location 0 contain the name of the user's job coded in EBCDIC. The two half words starting at locations 8 and A point to the user regions immediately following and preceding this region, thus defining the space addressable by this job. A user region can be a job, a symbiont, or a free region. Location C contains the number of bytes assigned to this job.

Location 120 contains the switch priority, if a switch command is entered that is to be in effect for the entire job. Otherwise, this byte will be zero.



Location 123 through 127 will contain the low-order five bytes of the current PSW (containing the interrupt code and instruction address) when a job abnormally terminates. Job control will write this information to the job log and display.

Location C0 contains the address of the active file table, which is used as an interface between the job and the files assigned to it (see 4.4.2).

### 2.2.3. Symbionts

Symbionts are dynamic extensions of the supervisor which occupy user main storage in much the same way as do user jobs. Each symbiont has a job-like preamble, TCBs preceding the code executed in the symbiont, and other necessary structures.

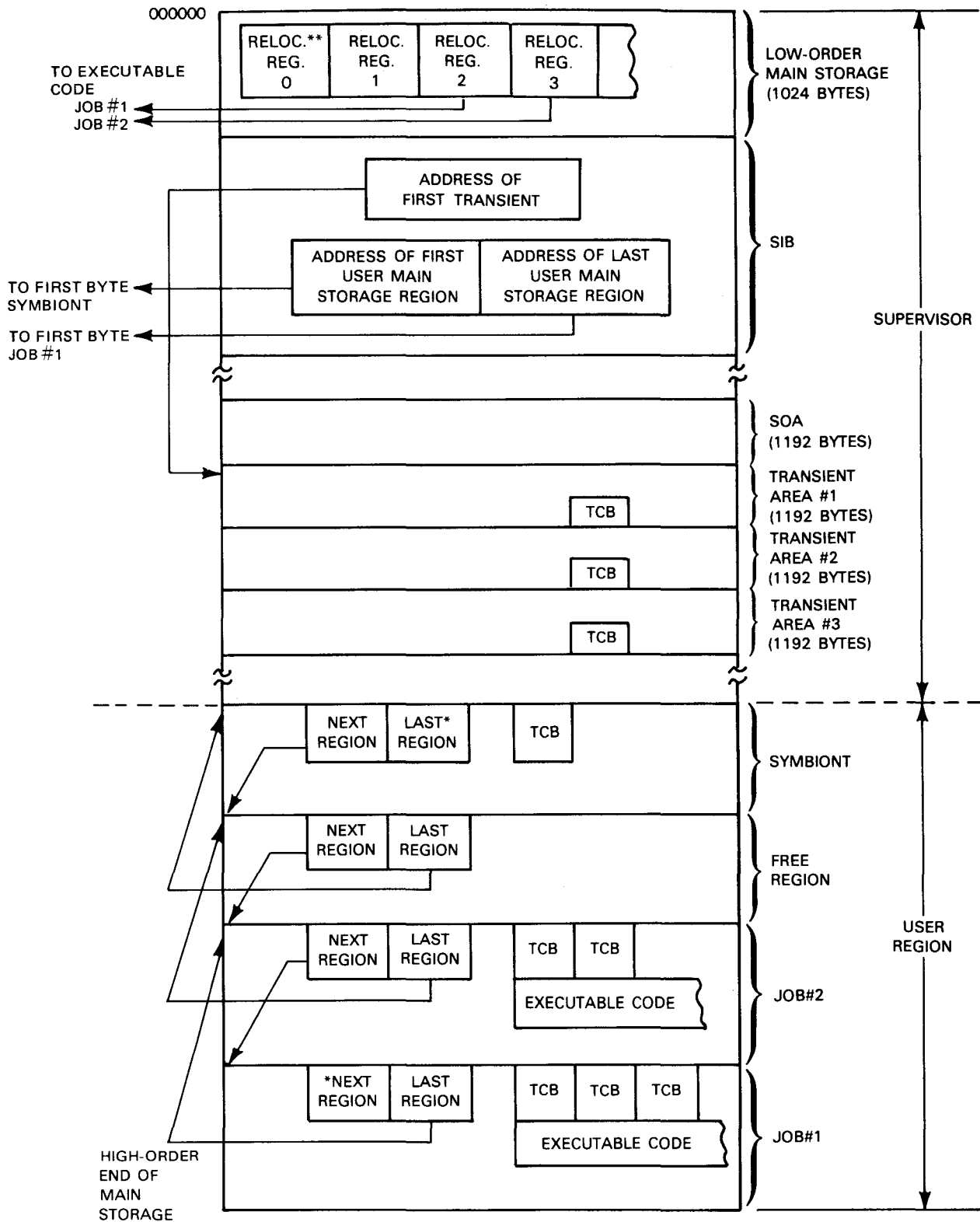
Some of the differences between symbionts and jobs lie in priority (see 2.3.1) and relocation keys: symbionts use key 0 (the supervisor) and run in supervisor mode; also, symbionts are loaded from the low-order end of user main storage upwards and are scheduled and loaded separately from job control.

### 2.2.4. Free Regions

As explained in 2.2, free regions are those parts of the user region not allocated to user jobs or symbionts. These regions are not simply ignored in OS/3, however: like a job or symbiont, each such region has a preamble with pointers in locations 8-B indicating the jobs/symbionts immediately preceding and following it, thus maintaining forward and backward links throughout the entire user region.

### 2.2.5. Structure Summary

The separate structures and pointers of an OS/3 system can be pulled together into a single diagram, shown in Figure 2-2.



\*Contains 0's indicating that it does not point to another region. \*\*Relocation registers not in main storage.

Figure 2-2. Summary of OS/3 in Main Storage

In Figure 2-2, two jobs and one symbiont are running concurrently, leaving a free region between them. Note that two relocation registers in low-order main storage are each pointing to the first byte of executable code in a job region. In analyzing dumps caused by errors within a user job, these relocation values can help determine specifically what job caused the error. Note that each job can have more than one TCB, although only one TCB can be active at any time.

## 2.3. Functional Description

This subsection summarizes some of the major operations of OS/3 software. You can find more information in the *Supervisor Technical Overview* (UP-8831). For dump analysis purposes, this discussion focuses on OS/3 tasks and task handling.

### 2.3.1. TCBs and Multitasking

As described previously, OS/3 tasks are processes that compete with other tasks for processor time. Each task is defined by an associated TCB. Each job or symbiont in the system is considered a primary task to which other tasks may be attached. The supervisor overlay and transient areas are individual tasks having one TCB each.

Tasks can be active or waited. An active task is one capable of getting processor control, but a waited task is made ineligible for processor control and must wait, usually for some event to occur, until it can be made active again.

Tasks are allotted processor time by the task switcher, a supervisor routine that scans all active tasks and chooses one among them to get the processor's attention. The TCB having the highest priority gets the processor first. Figure 2-3 shows the relative priorities assigned to different classes of TCBs.

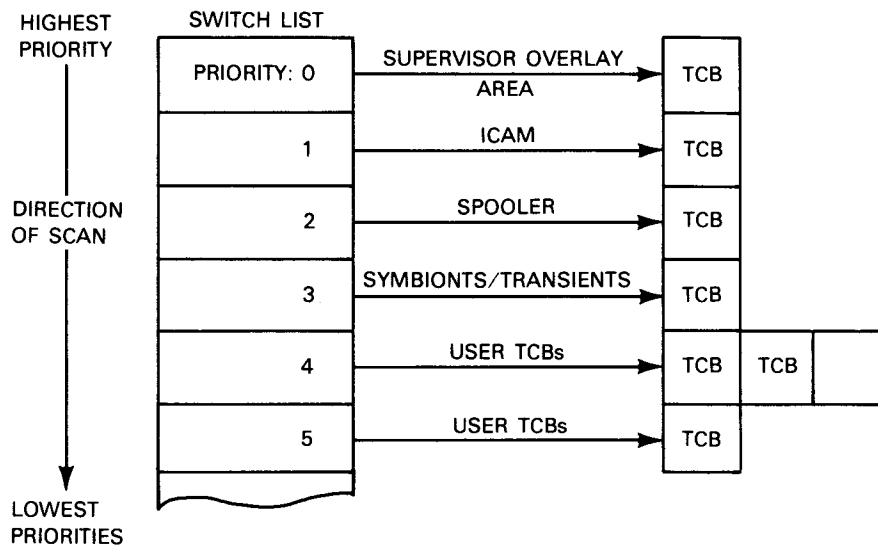


Figure 2-3. OS/3 Task Priorities

Note that the supervisor overlay TCB has the highest priority since supervisor routines must be run as soon as they are needed. As you move down the list, the tasks move from supervisor-oriented to user-oriented until you reach priority 4, the highest priority a user task can have.

Some priority levels have more than one task, although the tasks may be associated with different jobs and routines. When one task finishes or is interrupted, the switcher usually gives processor time to the next TCB at that priority if one exists and is active. Otherwise, the switcher scans lower priorities until it finds another active TCB to which to give processor control.

### 2.3.2. Task Switching Considerations

Switcher activity within a system can be determined from a dump and can thus help tell you why an error occurred within a job (which, remember, is itself a task operated on by the switcher).

Much task switching is done in response to an interrupt. When an interrupt is processed (Figure 2-4a) the current PSW is stored in the appropriate old PSW region of low-order main storage (step 1A). The corresponding new PSW is then loaded into the current PSW, thus handling processor control to an interrupt handling routine (step 1B). That routine in turn copies the old PSW into the double word PSW save area in bytes 20-27 of the interrupted task's TCB (step 2). Next, the routine in control stores the 16 registers used by the interrupted task into bytes 28-67 of the same TCB (step 3). (This step is omitted if control goes back to the same task as the one interrupted.) Finally, the interrupt handler returns processor control to the switcher.

The result of these operations is that control can be easily returned to the interrupted task when the time comes. This is shown in Figure 2-4b when the switcher gives control to a task which, like the task in Figure 2-4a, has had its registers and PSW saved. The switcher simply loads the 16 registers from the register save area (step 1) and the current PSW from the PSW save area (step 2). Control then passes to the instruction pointed to by the now-current PSW and the newly activated task begins processing.

As you will see later, task switching within a job is often prompted by interrupts responding to errors: in these cases, the old PSW (saved in the TCB of the interrupted task) shows which instruction caused the error. This is one of the most important techniques of dump analysis and is used extensively in the dump analyses in Section 4.

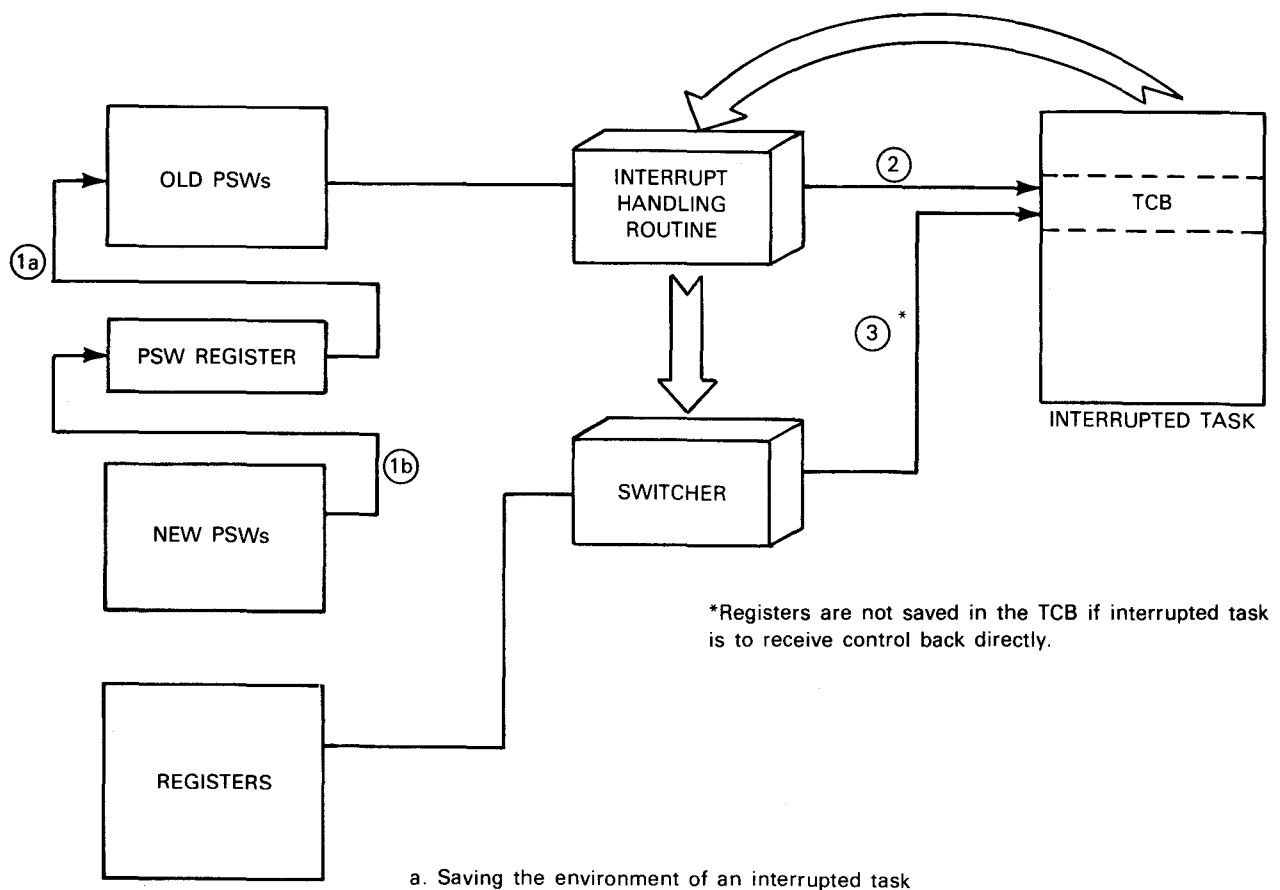


Figure 2-4. Passing Control from Task to Task (Part 1 of 2)

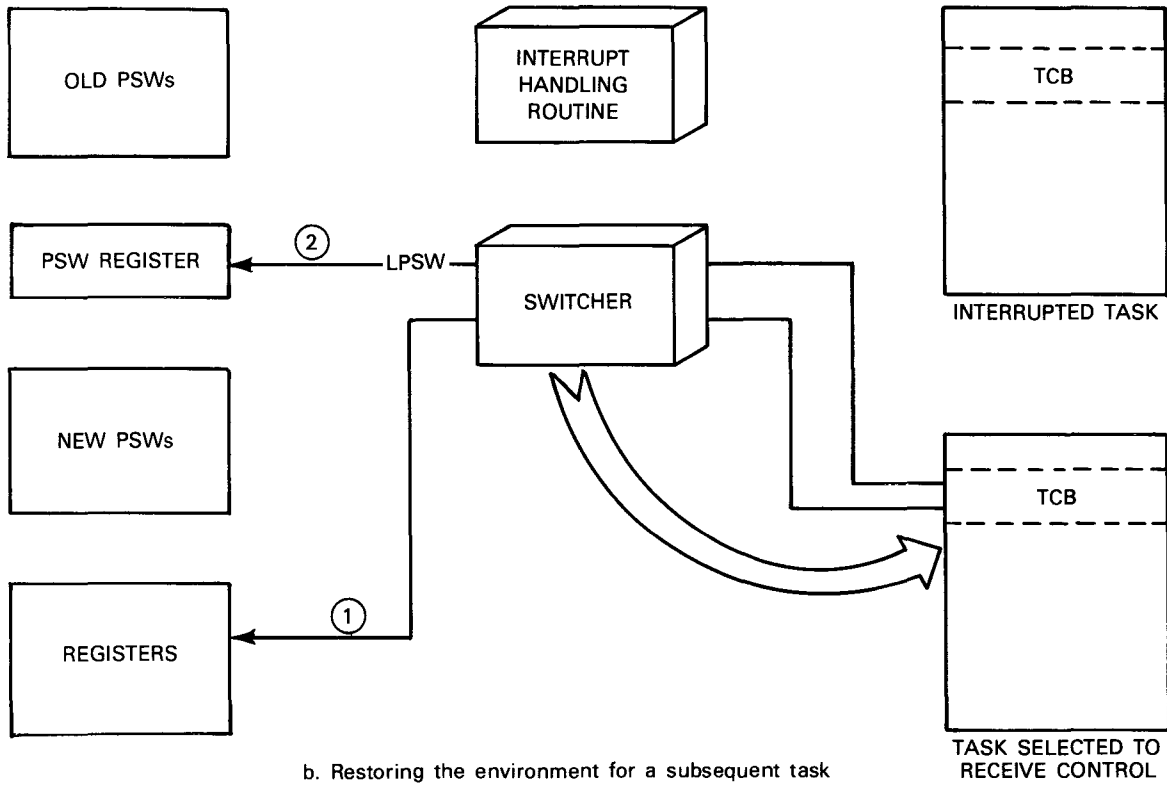


Figure 2-4. Passing Control from Task to Task (Part 2 of 2)

# Section 3

## Dumps and Their Formats

### 3.1. General

OS/3 provides three types of dumps. While all are basically hexadecimal images of main storage at the time they are taken, they differ in scope and purpose as follows:

- SYSDUMP

Dumps all or part of main storage and is run in two phases: main storage write and dump printout. In addition to the raw dump, the printout can provide a picture of your system in charts and text.

To produce an OPTION SYSDUMP, the RUN pack must have enough contiguous free space to store a copy of the system's main storage. See Appendix B for a procedure to allocate enough file space for your system.

- JOBDUMP

Dumps a user's job region upon abnormal termination of the job or execution of the DUMP or CANCEL macros. The raw dump is supplemented by charts and text interpreting the state of the job. JOBDUMP requires space on the first VTOC cylinder of SYSRUN to allocate files. You should eliminate unused files on SYSRUN.

- EOJ Dump

Dumps a user's job region without translation as well as the registers and the PSW.

The three types of dumps contain some common features, particularly in the user regions. The type of dump you choose to take will depend on some of the following considerations:

- HALT AND PROCEED (HPR) errors often require a SYSDUMP so that Unisys representatives can analyze system problems.
- System loops or hangs, usually appearing as operator communications that receive no reply from the system console, require a SYSDUMP. Like HPR errors, these usually require that a Unisys representative analyze your dump for system problems.

- User job errors generally occur within the user's own job region, thus requiring nothing more than a JOBDUMP or EOJ dump. These are the dumps you are most likely to work with in analyzing user program errors.

## 3.2. System Dump Routine (SYSDUMP)

### 3.2.1. SYSDUMP Makeup

The system dump routine (SYSDUMP) is most often used to determine why your system failed. But it may be used at any other time (even during normal processing). For example, you can use it interactively to dump all or specific regions of your job and associated system activity. SYSDUMP provides you with a listing that translates the state of the operating system when the dump was taken. In the case of a system failure, the operating system is shown just as it was at the time of the failure.

SYSDUMP is made up of two parts:

- A translated dump, which translates the state of the entire operating system into charts and text
- A labeled hexadecimal/character main storage dump

SYSDUMP is part of the supervisor and is always included at system generation (SYSGEN) time. It is designed to be run in a multiprogramming environment.

To get a system dump, the following two phases must occur:

1. The supervisor must write an image of main storage to the system dump file (\$Y\$DUMP). This phase is called the main storage write phase.
2. The SYSDMP load module, residing in load library \$Y\$LOD, must be called in and run. It is SYSDMP that reads the \$Y\$DUMP file and prints the translated and main storage dumps. This phase is called the execution phase.

**Note:** *The system dump file, \$Y\$DUMP, does not have to be on the SYSRES disk. It can be assigned to any disk. Refer to the Installation Guide (UP-8839) for more information.*

Generally, the execution phase of SYSDUMP immediately follows the main storage write phase. (The only exception to this is a SYSDUMP obtained from the operator controls. For this the operator must call SYSDMP; see 3.2.2.) Because the contents of main storage at the time of the main storage write are frozen in \$Y\$DUMP, you need not worry about the execution phase or any other job changing the contents of the system dump before you see the listing. This also means that you may use the SYSDUMP SAVE and RESTORE options (see 3.2.4) to perform a main storage write phase at one site and an execution phase at an entirely different site.



The operating system can generate a system dump in response to action by the operator, in response to a system error, or upon encountering an error in your job.

### 3.2.2. Operator-Initiated SYSDUMP

You can initiate a dump from either the operator controls or the system console/console workstation. Perform Step 1A if you are using the operator controls (located on the console workstation); perform Step 1B if you are using the system console.

**Note:** *The handbook referenced in these procedures is the Operations Guide (UP-8859).*

#### Step 1A. Starting from the operator controls

To initiate the main storage write phase, follow this procedure:

- For models 3 through 6:

While pressing the FUNCTION key:

- Press the D key to place the system in debug mode.
- Press the RESTART key.

The dump is completed when the following appears on the screen:

IPL

- For models 8 through 20:
  - Press the ESCAPE key on the console, then press M (for maintenance). A menu appears on the screen.
  - Select L (for system reset) in the menu and transmit. The same screen is displayed.
  - Press U (for run), then transmit. (Alternatively, you may press ESCAPE, followed by R, a system function.) When the dump is successfully written to the \$Y\$DMP file on disk, the screen shows 99999999 in the bottom right corner of the screen.
  - Perform an initial program load (IPL) on the system according to the directions in the *Operations Guide* (UP-8859).

When the IPL operation is finished, SYSDUMPO (the job stream that runs the SYSDMP module) will run automatically.

**Note:** *After Step 1A, you may get an HPR of 999F. This means that the \$YSDUMP file still contains information to be processed by SYSDMP (from a previous main storage write operation that has not yet been printed). If this happens, you can:*

- *Print it.* - Ordinarily, the information in \$YSDUMP is the system dump data most helpful to the Unisys representative who maintains your system. To print it, perform another IPL on the system according to the procedure in Step 1A.
- *Overwrite it.* - If you are sure you don't need the information in \$YSDUMP, you can write over it with an image of the system's main storage at the time you perform Step 1A. To do this, press
  - The FUNCTION and START console keys for models 3 through 6
  - The letter U and transmit (or the ESCAPE key followed by the letter R) for models 8 through 20

*and wait for memory write completion. Perform an IPL on the system according to the directions in the Operations Guide (UP-8859). When the IPL operation is finished, the SYSDUMPO jobstream will run automatically.*

### Step 1B. Starting from the system console

To initiate a dump, key in the following:

```
SYSDUMP
```

The SYSDUMP command writes main storage to the \$YSDUMP file, then enters a job in the job queue, whose only function is to print the system dump. The job is named SYSDMPxx, where xx is the unique number assigned to this job by the system. When the job is scheduled and run, the following message appears on the system console:

```
DUMP OPTION (ALL, NONE, DUMP, JOBS, EDIT, MINI, SAVE, RESTORE)
```

See Step 2 for an explanation of these parameters.

**Note:** *After the supervisor writes main storage to \$YSDUMP, it locks that file until job SYSDMPxx is complete. If job SYSDMPxx is removed from the system prior to displaying the DUMP OPTION message, you must unlock the \$YSDUMP file by entering the following console command:*

```
SET SY,LOFF
```

Two examples of when this command must be used are the following:

1. If you delete job SYSDMPxx from the job queue before it is scheduled
2. If the run processor encounters an error while trying to put job SYSDMPxx in the job queue

## Step 2.

Enter the SYSDUMP command in one of the following three formats, depending on the output you want (unless the supervisor has already called it automatically).

### Format 1

```
RV SYSDUMP
```

Produces an output identical to that produced by running JOB SYSDUMPO with the dump option ALL. (See Formats 2 and 3).

### Format 2

$$RV \text{ SYSDUMPO } \left[ \begin{array}{l} \text{,,DO=} \left( \begin{array}{l} \underline{\text{ALL}} \\ \underline{\text{NONE}} \\ \underline{\text{DUMP}} \\ \underline{\text{JOBS}} \\ \underline{\text{EDIT}} \\ \underline{\text{MINI}} \\ \underline{\text{SAVE}} \\ \underline{\text{RESTORE}} \end{array} \right) \left[ \begin{array}{l} \text{,V=} \left\{ \begin{array}{l} \text{vsn} \\ \text{(vsn[,A],[dev])} \end{array} \right\} \end{array} \right] \left[ \text{,P=did} \right] \end{array} \right]$$

### Keyword Parameter DO

$$\text{DO=} \left( \begin{array}{l} \underline{\text{ALL}} \\ \underline{\text{NONE}} \\ \underline{\text{DUMP}} \\ \underline{\text{JOBS}} \\ \underline{\text{EDIT}} \\ \underline{\text{MINI}} \\ \underline{\text{SAVE}} \\ \underline{\text{RESTORE}} \end{array} \right)$$

Specifies the dump option you want. Available options are:

ALL

The dump listing displays and translates the state of your entire system, including a header page, low-order main storage, the system information block, the physical unit block, the main storage map (giving the locations of everything in main storage), the system switch list, the job region for each job or symbiont in the system, and the free region. Essentially, the ALL option combines the DUMP and EDIT options, supplying a hexadecimal dump with an English translation.

NONE

No output listing is produced.

DUMP

Writes the entire main storage in hexadecimal format.

JOBS

Produces a hexadecimal dump and a translation for all the jobs, symbionts, and shared code in main storage at the time it was written to the \$Y\$DUMP file.

EDIT

Provides an English language description of the state of the system.

*Note: EDIT replaces TRANSLATED. Both are supported and they are functionally equivalent.*

MINI

Provides a printed hexadecimal dump, including a table of contents, of specific main storage regions dependent upon the HPR or system error code. This is the recommended option for dumps to be sent to Unisys for analysis.

SAVE

Saves the \$Y\$DUMP file to magnetic tape or diskette (see 3.2.4).

RESTORE

Restores the \$Y\$DUMP file from magnetic tape or diskette (see 3.2.4).

If you omit keyword parameter DO or if the supervisor has already called SYSDMP (as with the SYSDUMP console command), then the system will display the following message:

```
DUMP OPTION (ALL, NONE, DUMP, JOBS, EDIT, MINI, SAVE, RESTORE)
```

You then enter the *one* option you want. SYSDMP redisplays this message after finishing a SAVE or RESTORE function, so you can perform other SYSDMP operations within the same job.

**Keyword Parameter V**
$$V = \left\{ \begin{array}{l} \text{vsn} \\ (\text{vsn}[,A],[\text{dev}]) \end{array} \right\}$$

This parameter allows you to display a dump not residing on the booted SYSRES device.

**vsn**

If specified, all data to be processed reside on the specified vsn. If not specified, the booted SYSRES is used.

**A**

Indicates that the SYSDMP load module will also be executed from the specified vsn. The vsn item must be an alternate SYSRES disk.

**dev**

Indicates that the device number of the disk pack is specified. Use this option when you have two volume serial numbers that are the same.

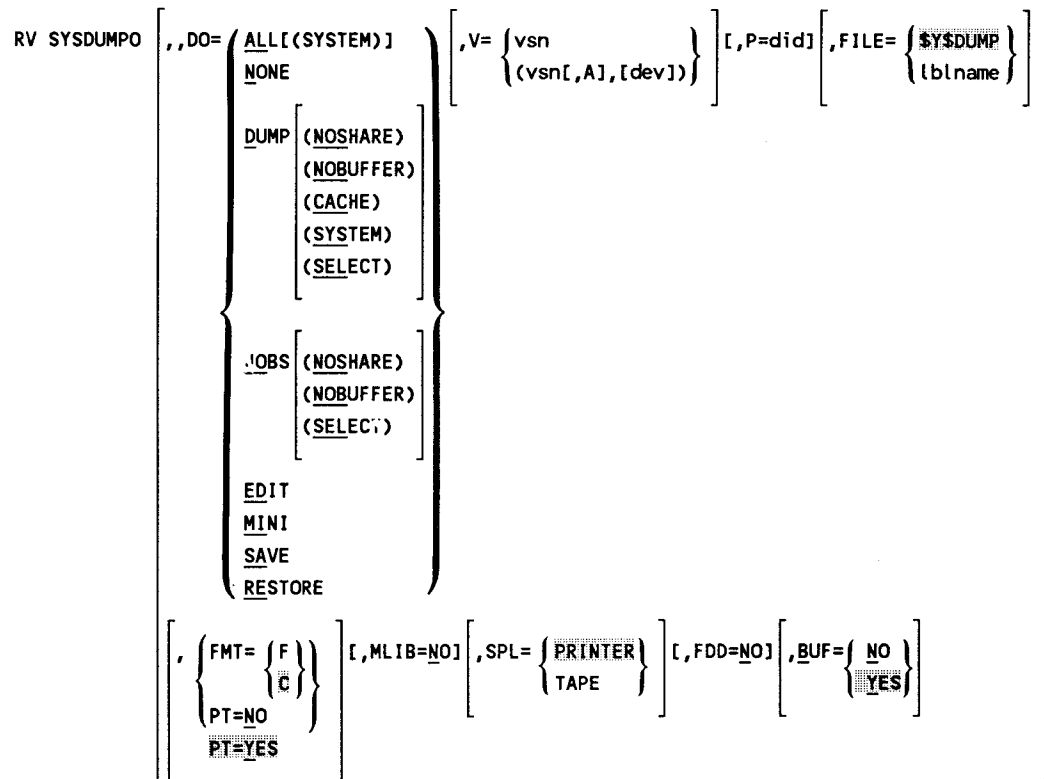
**Keyword Parameter P**

**P=did**

If you wish to assign a specific printer as the output device for SYSDUMP, enter its 3-character device address using this keyword parameter. You may wish to do this if your system supports spooling but you do not want to spool the SYSDUMP listing.

If you want to enter more specifics in the RV SYSDUMPO command, use the suboptions shown in Format 3.

**Format 3**



**Note:** Of the Format 3 suboptions, *SELECT* is mutually exclusive of *NOSHARE*, *NOBUFFER*, *CACHE*, and *SYSTEM*.

**Keyword Parameter DO**

Same as in Format 2, except with the following suboptions and options added.

**(NOSHARE)**

Excludes shared code modules in SYSDUMPs. *NOSHARE* can be used with any SYSDUMP option except with *SAVE*, *RESTORE*, or *SELECT*.

**(NOBUFFER)**

Excludes the display of dynamic buffers and the dynamic shared code.

**(CACHE)**

Displays the entire *CACHE* main storage. If not selected, dump options *JOB* and *DUMP* will display only 6K of *CACHE*.

**(SYSTEM)**

Displays the following preselected regions of the supervisor

Low Core

System Information Block (SIB)

Switch List

Physical Unit Blocks Table (PUBS)

Shared Code Directory

Supervisor Debug Tables

Supervisor Overlay Area (SOA)/Transient Regions

System Spool Control Table

Dynamic Buffer Information Block

Resident Buffer Control Block

Shared Code Information Block

**(SELECT)**

Allows you to interactively select particular regions of the main storage dump. *This suboption creates an interactive environment at the operator's console. You will then be prompted to enter requests as follows:*

```
SD28  SELECT(DISPLAY, JOBNAME, MEMORY, REGION, SHARED...  
SD28  ....TRANAREA, TERMINATE)
```

An abbreviated version of this message:

```
SD28  SELECT(DIS, JOB, MEM, REG, SHAR, TRAN, TERM)
```

appears after the first selection is made. Thereafter, the complete message is displayed only after an incorrect response or when the HELP prompter is invoked.

**Note:** *At any time during entry of SELECT suboptions, you can enter a SELECT keyword followed by a blank space, and a help prompter for that keyword will appear on screen.*

You can then enter one of the following keyword requests:

$$\left( \begin{array}{l} \text{DISPLAY=} \left( \begin{array}{l} \text{startaddr:nn} \\ \left\{ \begin{array}{l} \text{X'xxxxxxxx'} \\ \text{C'ccccccc'} \end{array} \right\} [\text{startaddr-endaddr}] \\ \left( \begin{array}{l} \text{CR} \\ \text{IO} \\ \text{PR} \\ \text{RR} \\ \text{SN} \\ \text{SR} \\ \text{JOBS} \end{array} \right) \end{array} \right) \\ \text{JOBNAME=} \\ \text{MEMORY=} \\ \text{SHARED=} \\ \text{TRANAREA=} \\ \text{REGION=} \\ \text{TERMINATE=} \end{array} \right)$$

where:

DISPLAY=startaddr:nn

Displays 1 to 10 lines of main storage on the system console or 1 to 20 lines on a user workstation, from the start address forward. Enter a hexadecimal value for startaddr.

When this keyword is in effect, a blank response results in the next logical main storage addresses being displayed. The line count is the same as the previous response.

DISPLAY=  $\left\{ \begin{array}{l} \text{X'xxxxxxxx'} \\ \text{C'ccccccc'} \end{array} \right\} [(\text{startaddr-endaddr})]$

Displays the address of the first occurrence of the sequence specified. For the address of the next occurrence, return a blank response to the SELECT request with this keyword in effect.

where:

xxxxxxxx

Is a hexadecimal character string of up to 16 bytes.

ccccccc

Is an ASCII character string of up to 16 characters.



startaddr-endaddr

Limits the search for the hexadecimal or ASCII character string to the range specified.

DISPLAY= ( CR  
IO  
PR  
RR  
SN  
SR  
JOBS )

Displays the contents of specified registers.

where:

CR

Represents the control registers.

IO

Represents the I/O relocation registers (models 3 through 6 only).

PR

Represents the problem registers.

RR

Represents the relocation registers.

SN

Represents the snap registers. (Not available on models 3 through 6.)

SR

Represents the supervisor registers.

JOBS

Lists the names and addresses of all the jobs in the system.

JOBNAME=job/symbiont-name

Specifies an 8-character job/symbiont name, which produces a hexadecimal dump of that job/symbiont. If you choose the JOB option, an English language description is also provided.

If there is more than one copy of a job or symbiont, SYSDUMPO displays a list of the duplicate jobs or symbionts and their addresses. You can then choose the copy you want. For example:

```

15?SYSDMP76 SD28   SELECT (DIS, JOB, MEM, REG, SHAR, TRAN, TERM)
15 J=SL$$OWOO
16 SYSDMP76 SD40   THERE ARE 002 COPIES OF SL$$OWOO
17 SYSDMP76 SD40   ENTER NUMBER OF DESIRED ADDRESS. (0=NONE)
18?SYSDMP76 SD40   ENTER 1=06D100 2=076800
18 1
19?SYSDMP76 SD40   SELECT (DIS, JOB, MEM, REG, SHAR, TRAN, TERM)
    
```

$$\text{MEMORY} = \left( \begin{array}{l} \text{startaddr} \text{ -endaddr} \\ \text{ALL} \end{array} \left( \begin{array}{l} \text{-endaddr} \\ \text{,byte-count} \\ \text{[-END]} \\ \text{,END} \end{array} \right) \right)$$

where:

startaddr-endaddr

Dumps main storage from the start address to the end address specified. Note the use of the dash. Enter hexadecimal values for startaddr and endaddr.

startaddr,byte-count

Dumps the number of bytes of main storage specified, starting at the address specified. Note the use of the comma. Enter a hexadecimal value for startaddr and a decimal value for byte-count.

startaddr-END

Dumps main storage from the start address to the end of machine capacity. Enter a hexadecimal value for startaddr.

startaddr,END

Same as startaddr-END.

ALL

Dumps all the main storage.

SHARED=module-name

Specifies an 8-character shared code module name whose module is dumped.

$$\text{TRANAREA} = \left( \begin{array}{l} \text{ALL} \\ \text{decimal-area-number} \\ \text{D'ID'} \\ \text{X'ID'} \end{array} \right)$$

where:

ALL

Dumps all transient areas.

decimal-area-number

Dumps a specified transient area.

D'ID'

Decimal ID of the requested transient area.

X'ID'

Hexadecimal ID of the requested transient area.

REGION

You may enter one of the following acronyms or words designating a particular region:

$$\left( \begin{array}{l} \text{CACHE} \\ \text{DBIB} \\ \text{DEBUG} \\ \text{DISABLED} \\ \text{FREE} \\ \text{ISIB} \\ \text{LOWCORE} \\ \text{PUBS} \\ \text{RBCB} \\ \text{SCD} \\ \text{SCIB} \\ \text{SIB} \\ \text{SOA} \\ \text{SUPERVISOR} \\ \text{SWITCHLIST} \\ \text{TRACE} \\ \text{UNIDENTIFIED} \end{array} \right)$$

where:

CACHE

Cache module and buffers

DBIB

Dynamic buffer information block

DEBUG

System debug tables

DISABLED

Main storage as disabled

FREE

Main storage as free

ISIB

Interactive services information block

LOWCORE

Main storage from location zero to system information block (SIB)

PUBS

Physical unit blocks table

RBCB

Resident buffer control block

SCD

Shared code directory

SCIB

Shared code information block

SIB

System information block

SOA

Supervisor overlay area

SUPERVISOR

Dumps the entire supervisor region

SWITCHLIST

Switch list

TRACE

System trace table

UNIDENTIFIED

Unlabeled region of main storage (MEMORY) in synopsis

**TERMINATE**

Ends the dump session. When a selection has been processed, you will always be prompted to enter another selection until you choose **TERMINATE**. A keyword remains in effect until you choose **TERMINATE**. For example, if you select **REGION=SIB**, you may then enter any accepted region acronym without repeating the **REGION=keyword**.

**Keyword Parameter V**

The parameters are defined in Format 2.

**Keyword Parameter P**

The parameters are defined in Format 2.

**Keyword Parameter FILE**
$$\text{FILE} = \left\{ \begin{array}{l} \text{\$Y\$DUMP} \\ \text{lblname} \end{array} \right\}$$

where:

**FILE=\$Y\$DUMP**

This default parameter allows a dump image to be processed from system file **\$Y\$DUMP**.

**FILE=lblname**

Allows a dump image to be processed from some file other than system file **\$Y\$DUMP**.

**Keyword Parameter FMT**
$$\text{FMT} = \left\{ \begin{array}{l} \text{F} \\ \text{C} \end{array} \right\}$$

where:

**FMT=F**

Provides a full System Maintenance Change (SMC) listing after the dump.

**FMT=C**

Provides a condensed SMC listing after the dump. (This is the default.)

### Keyword Parameter PT

$\left\{ \begin{array}{l} \text{PT=NO} \\ \text{PT=YES} \end{array} \right\}$

where:

**PT=NO**  
System Maintenance Change (SMC) listing is not printed. If this is entered, you cannot specify any FMT.

**PT=YES**  
SMC listing is printed. If nothing else is specified (FMT or PT=NO), this is the default.

### Keyword Parameter MLIB

**MLIB=NO**  
Prevents dumping of the \$Y\$SDF file.

### Keyword Parameter SPL

$\text{SPL} = \left\{ \begin{array}{l} \text{PRINTER} \\ \text{TAPE} \end{array} \right\}$

where:

**SPL=PRINTER**  
Sends all spooled output to the printer. (This is the default.)

**SPL=TAPE**  
Sends all spooled output to tape. If this is not specified, then output is printed.

### Keyword Parameter FDD

**FDD=NO**  
Turns off printing of the FDDO diskette.

**BUF=NO**  
Turns off printing of the dynamic buffer summary.

**BUF=YES**  
Turns on printing of the dynamic buffer summary. (This is the default.)

Table 3-1 is a summary of the run statement keywords.

Table 3-1. Run Statement Keyword

Keyword Parameter	Option	Suboption	Format
DO	<u>ALL</u>		DO=AL
		<u>SYSTEM</u>	DO=A(SYS)
	<u>NONE</u>		DO=NO
	<u>DUMP</u>		DO=DU
		<u>NOSHARE</u>	DO=DU(NOS)
		<u>NOBUFFER</u>	DO=DU(NOBU)
		<u>CACHE</u>	DO=DU(CAC)
		<u>SYSTEM</u>	DO=DU(SYS)
		<u>SELECT</u>	DO=DU(SEL)
		<u>JOBS</u>	DO=JO
		<u>NOSHARE</u>	DO=JO(NOS)
		<u>NOBUFFER</u>	DO=JO(NOBU)
		<u>SELECT</u>	DO=JO(SEL)
		<u>EDIT</u>	DO=ED
		<u>MINI</u>	DO=MI
		<u>SAVE</u>	DO=SA
	<u>RESTORE</u>	DO=RE	
V	vsn		V=vsn
	(vsn,A)		V=(vsn[,A],[dev])
P	did		P=did

continued

Table 3-1. Run Statement Keywords (cont.)

Keyword Parameter	Option	Suboption	Format
FILE	SYSDUMP <sup>1</sup>		FILE=SYSDUMP
	lbname		FILE=lbname
FMT	F		FMT=F
	C <sup>1</sup>		FMT=C
PT	<u>N</u> O		PT=N
	<u>Y</u> ES <sup>1</sup>		PT=Y
MLIB	<u>N</u> O		MLIB=N
SPL	PRINTER <sup>1</sup>		SPL=PRINTER
	TAPE		SPL=TAPE
FDD	<u>N</u> O		FDD=N
BUF	<u>N</u> O		BUF=N

<sup>1</sup> Default

### 3.2.3. Supervisor-Initiated SYSDUMP

The supervisor will automatically generate a system dump in either of two situations:

- An error occurs within a user program.
- An error occurs within the supervisor itself.

These two situations are discussed in more detail in 3.2.4.

#### SYSDUMP from User Program Errors

You can enable the supervisor to generate a SYSDUMP if it encounters a program error or the program executes either the DUMP or CANCEL macroinstruction. In this case, the main storage write phase and execution phases are called and executed automatically, one after the other. To accomplish this, the supervisor calls and executes the SYSDMP load module in your job region within the same step in which the program error occurred.



To enable a supervisor-initiated SYSDUMP you must:

- Include a printer device assignment set with the LFD name PRNTR
- Include //ΔOPTION SYSDUMP job control statement preceding the //ΔEXEC statement of the program for which the dump is to be taken

If you want to enable the SYSDUMP option for more than one step, use the // OPTION GSYSDUMP job control statement. Inserting it in your control stream enables the SYSDUMP option for all steps following it right up to the end of the job.

If you want to restart a job while OPTION SYSDUMP is still processing, use the // OPTION PSYSDUMP job control statement. This option immediately terminates the failing job and processes the OPTION PSYSDUMP under a cover name of JOBNAME#.

If a program run with //ΔOPTION SYSDUMP fails, the supervisor will write a main storage image to \$Y\$DUMP, load the SYSDMP module for the execution phase, then display the following message:

```
DUMP OPTION (ALL, NONE, DUMP, TRANSLATED, JOBS, RESTORE, SAVE)
```

You may want to tell the system operator beforehand what SYSDUMP option to enter should a system dump be generated from your program. See Step 2 in 3.2.2 for an explanation of these options.

## SYSDUMP from System Errors

Some system errors (errors occurring within the supervisor) may automatically initiate a SYSDUMP. The supervisor will attempt to write a main storage image to \$Y\$DUMP; if it succeeds, an SE15 message will be displayed on the system console, as in the following example:

```
SE15 SYSTEM ERROR 20 IN TRANS # 33-SYSDUMP WRITTEN TO DISK
```

The system will then schedule job SYSDMPxx to print the contents of \$Y\$DUMP. This example of message SE15 shows that a program check (error code 20) occurred in transient number 33. The SYSDUMP WRITTEN TO DISK message indicates that job SYSDMPxx has been scheduled. When you see the DUMP OPTION message, you should reply with the DUMP, MINI, or SAVE option, and send the resulting printout to your Unisys representative for analysis.

If the `Y$DUMP` file is in use or locked when a system error occurs, a SE16 message will be displayed on the console showing the error and the current contents of the program status word (PSW) and registers. In this case, no main storage write will be performed. The system will, however, continue to run.

### 3.2.4. SYSDUMP SAVE and RESTORE Options

The end product of a SYSDUMP is the printed listing that shows the state of the operating system and user jobs at the time the system crashed. For analysis purposes you may wish to send a SYSDUMP listing to another site. The size and bulk of a printed listing, however, may make sending it a slow and expensive process. To help you avoid this problem, Unisys provides the SAVE and RESTORE options with SYSDUMP. You may use the SAVE option to copy the `Y$DUMP` file to diskettes or a magnetic tape volume. With the RESTORE option you copy the data on that diskette or tape back onto the `Y$DUMP` file. This means, for example, that at one site you can copy a `Y$DUMP` file to a diskette (with SAVE), send only the diskette to another site, copy the diskette data to the `Y$DUMP` file there (with RESTORE), and run SYSDUMPO to get the printed listing of the dump, all at reduced cost and in less time.

You can use the SAVE and RESTORE options with single and multivolume tapes and diskettes; however, there are some requirements.

- Tape
  - Requires Consolidated Data Management (CDM)
  - Requires 9-track tape
- Diskette
  - Requires Consolidated Data Management

#### SAVE Option

To use the SAVE option you must first prep a diskette or magnetic tape using the canned job control stream SD\$PREP. When you call SYSDUMPO and get the DUMP OPTION message, you key in SAVE. The system will reply with the following console message:

```
ENTER DUMP DEVICE INFORMATION (DISKETTE, TAPE, NONE)
```

Key in the device you will use in saving the \$YSDUMP file, DISKETTE or TAPE, or key in NONE to terminate SAVE without saving the \$YSDUMP file. If you enter DISKETTE or TAPE, the system will then ask you to mount a prepped diskette or magnetic tape on an unused drive. When you have done so, the system will copy the \$YSDUMP file. After it finishes it will display:

```
$YSDUMP SAVED ON TO { TAPE }
                   { DISKETTE }
```

The system will then display the DUMP OPTION message so that you may select another SYSDUMPO option or terminate SYSDUMPO altogether (with NONE). In other words, when you run SYSDUMPO with the SAVE option, you will be assigning a diskette/tape in addition to supplying a brief "problem description".

You can send tapes containing saved dumps to the Customer Support Center for analysis. Please include the following on the label of the tape:

- Whether or not the system that created the tape was generated for tape block numbering
- Streaming or non-streaming type of the tape drive
- Your company name
- Applicable reference number
  - UCF number (supplied by Unisys)
  - Authorization number (supplied by Unisys)
- Date
- Number multivolume tapes giving the number of the tapes and the total number of tapes (for example: 1 of 3, 2 of 3, 3 of 3)
- The dump name (please note on UCF form in comments area), for example, HENRC001

*Note: The SYSDUMP will print the SMC LIST and error log. Include these listings with the tape or tapes.*

## RESTORE Option

You use the RESTORE option much like the SAVE option. In response to the SYSDUMPO DUMP OPTION message you key in RESTORE. The system will reply with the following console message:

```
ENTER DUMP DEVICE INFORMATION (DISKETTE, TAPE, NONE)
```

Key in the device type that contains the saved \$YSDUMP file, DISKETTE or TAPE, or key in NONE to terminate RESTORE without restoring the \$YSDUMP file. If you enter DISKETTE or TAPE, the system will then ask you to mount the magnetic tape or diskette. When you have done so, the system will copy the data to the \$YSDUMP on the SYSRES volume, destroying any data previously stored there. After it finishes, it will display:

```
$YSDUMP RESTORED FROM { TAPE
                       { DISKETTE }
```

The system will then display the DUMP OPTION message so that you may select another SYSDUMPO option or terminate SYSDUMPO altogether (with NONE). You can now have the system print the SYSDUMP listing in any of its available formats.

*Note: When the SYSDUMP main storage write and execution phases take place in the same system, the main storage map will include a list of the CSECTs in main storage. If you run SYSDUMPO on a system other than the one whose main storage image is in the \$YSDUMP file, you will not get the CSECT listing. All other parts of the SYSDUMP listing, however, will be printed.*

### Prepping a Diskette or Magnetic Tape for SAVE/RESTORE

Before you can use a diskette or magnetic tape with the SAVE or RESTORE options, you must prep it using the canned job control stream SD\$PREP. The control stream accepts a diskette or magnetic tape straight from the factory, formats it if necessary, and allocates a single-volume file called \$YSDMP, the file which the SAVE and RESTORE options use.

*Note: If you send a dump on tape to the Customer Support Center, make sure another tape is prepped and ready for future dumps.*

To call SD\$PREP, key in:

```
RV SD$PREP,,TYPE= { TAPE
                  { DISKETTE },NUMBER= { 1
                                       { nn },VOLUME= { 1
                                                       { vv },SIDES= { 1
                                                                           { 2 }
```

where:

```
TYPE= { TAPE
      { DISKETTE }
```

Specifies the type of device to be prepped.

```
NUMBER= { 1
        { nn }
```

nn specifies 01 to 16 tape or diskette volumes.

$$\text{VOLUME} = \left\{ \begin{array}{l} \text{1} \\ \text{vv} \end{array} \right\}$$

vv specifies the first volume to be prepped.

$$\text{SIDES} = \left\{ \begin{array}{l} \text{1} \\ \text{2} \end{array} \right\}$$

Specifies a 1- or 2-sided diskette.

The system will ask you to mount the diskette or magnetic tape to be prepped. When you have done so, the system preps it.

These are a few things to remember when using SD\$PREP:

1. Prepping a magnetic tape or diskette for SYSDUMPO destroys all data previously recorded on that medium.
2. A diskette or magnetic tape prepped by SD\$PREP cannot be used for other files or programs. To use a magnetic tape or diskette (for other files) that has already been prepped with SD\$PREP, you must prep it again using the appropriate system service program described in the *System Service Programs (SSP) Operating Guide* (UP-8841).
3. Tapes or diskettes prepped prior to Release 8 are no longer valid.

### 3.3. Job Dump Routine (JOB\_DUMP)

You use the job dump routines JOB\_DUMP and ABR\_DUMP to determine what caused your job to terminate abnormally.

#### 3.3.1. Full Job Dump (OPTION JOB\_DUMP)

JOB\_DUMP provides a method for determining what caused the job to terminate abnormally. It prints out a listing of the state of the job region when the job crashed. JOB\_DUMP is made up of two parts:

- A dump that translates the state of the job region into charts and text
- A labeled hexadecimal/character main storage dump

To execute `JOBDUMP`, a `// OPTION JOBDUMP` job control statement is required and must precede the `EXEC` job control statement. The first statement invokes `JOBDUMP` if the job terminates abnormally or upon execution of a `DUMP` or `CANCEL` macroinstruction in the assembler program.

There must also be a printer device assignment present in the control stream with a `LFD` file name of `PRNTR`.

### 3.3.2. Abbreviated Job Dump (OPTION ABRDUMP)

The abbreviated job dump provides you with a shortened listing of the full job dump to help you determine what caused the job to terminate abnormally. When `ABRDUMP` is called, only the area in the vicinity of the last instruction executed, along with the address and contents of the `I/O` buffers associated with the `OPEN DTFs`, is printed.

To execute `ABRDUMP`, a `// OPTION ABRDUMP` job control statement is used in place of the `// OPTION JOBDUMP` job control statement. The abbreviated job dump is called similarly to the full job dump.

## 3.4. EOJ Dump Routine

The `EOJ` dump routine is called by either the `DUMP` macro used in place of the `EOJ` macro in your assembler program or by an abnormal termination of your job. In either case, the `// OPTION DUMP` job control statement must be present in your control stream. The `EOJ` is in hexadecimal format and is divided into four sections: problem program registers, job preamble, task control blocks (`TCBs`), and your program region. In addition, the dump also gives you the `PSW` at interrupt time, the error code that caused the abnormal termination, and the next `TCB` address.

## 3.5. Supervisor Trace Analysis

The supervisor trace facility is automatically included in any system dump, whether operator- or supervisor-initiated; you don't need to activate it.

The module for the supervisor trace facility is `SM$TRACE`.

### 3.5.1. How to Read a Supervisor Trace

1. Find the module heading `SM$TRACE` in the system dump printout. It should be located near the end of the resident supervisor portion of the dump.
2. Find the word `TRAC` in the interpreted (edited) section on the right side of the system dump printout.

3. Now find the equivalent of TRAC in the hexadecimal portion of the printout. This hex value marks the beginning of the trace table where the supervisor activities are logged.

Note that the entry for each supervisor activity consists of four hexadecimal words. The supervisor trace table can contain up to 400 entries and will wrap (overwrite the beginning of the table) when the end of the table is reached, so that only the 400 most current entries are listed.

4. Find the word following the hex equivalent of TRAC. This word contains the address of the next available entry in the table (i.e., where the next entry will be logged).
5. Look for this address in the hex portion of the printout. Note the entry in *front* of this address. This entry is the last (most current) one logged; it is the logical end of the supervisor trace table.

Table 3-2 describes the supervisor area trace table entry formats. Note that word 1 for each entry appears in the interpreted portion at the right of the dump printout.

Table 3-2 also provides the interpretations of the hex values of words 2, 3, and 4 for each supervisor area. Note that all zeros indicate that there are no values (no entries to be made) for that word. In some cases, words 2 and 3 are combined to provide one value.

Table 3-2. Supervisor Area Trace Table Entry Formats

Supervisor Area	Word 1	Word 2	Word 3	Word 4
Timer	TI	00000000	00000000	Interval timer register
Machine check	MC	Machine check old program status word (PSW)		Interval timer register
Program check	PC	Program check old PSW		Interval timer register
Transient overlay	TO	Transient area address	Overlay transient id	Interval timer register
Task switch	TS	00000000	TCB address of new task	Interval timer register
SVC call	SV	SVC old PSW		Interval timer register
Transient release	TR	Transient area address	Transient id	Interval timer register
Shared code call	SC	TCB address of task using shared code	Name of shared code module being called	Interval timer register
Shared code release	SR	TCB address of task using shared code	Name of shared code module being exited	Interval timer register



# Section 4

## Sample Dump Analyses

### 4.1. General

In the preceding sections, you got a general overview of OS/3 dumps and of how OS/3 interacts with user programs. In this section, you will see examples of dump analyses applied to actual user programs by using the ideas already presented. A number of programs are shown here that have failed while being executed. Although all have been compiled or assembled correctly, errors have nonetheless remained. It is in uncovering these errors that dump analysis is a useful tool.

Three programs are presented in this section, each written in a language supported by OS/3:

- BALOBJ, a program that reads data from cards, adds the data, and outputs the result to a printer, is written in Basic Assembler Language (BAL).
- COBOBJ, a program that performs the same functions as BALOBJ, is written in COBOL.
- RPGOBJ, a program that performs the same functions as BALOBJ, is written in RPG II.

In these examples, you will be working with different types of dumps, as well as with other diagnostic tools supplied with the individual language processors. Each example takes the following form:

- An outline of the program and the circumstances under which it failed
- A list of the dumps and other materials used in the analysis
- A brief outline of the particular type of dump used or of an analysis technique
- A step-by-step dump analysis narrative
- An edited copy of the output listings used in the analysis

## Dump Analyses

The dump analysis itself will consist of two parts; the narrative and the listings. Most important among these listings is the dump itself, but compiler/assembler listings and other materials are also included where needed.

In a typical dump analysis, you will move forward and backward within and often outside the dump. With analysis, you are trying to uncover the logical (and not necessarily the physical) sequence of events leading up to a program failure.

To help you find your way through the following analyses, they are broken down into logical steps numbered ①, ②, ③, etc. With each step, a relevant portion of the listings is reproduced within which an identically numbered pointer indicates the data used in that step. All the pointers reappear in the listings themselves, together with an index indicating where each pointer is.

As an example, we show two excerpts from a dump analysis, one taken from the narrative and the other from the output listings.

### Dump Analysis Narrative

```
0001A0-00000000 00000000 00000000 00000000 00000000 00000000 F2F3F240 F3F6F540 *.....232 365 -001890
0001C0-000002C5 0303C5E8 68A0E3D6 04404640 40404040 40404040 40404040 40404040 * KELLEY, TOM -001880
0001E0-F1F0F340 40F3F0F2 47A0E5C9 05C3C5D5 E36846C7 C5D6D9C7 C5404040 40404040 *103 362 VINCENT, GEORGE -001800
```

At location 18AD ⑨ we see that the three bytes operated on by the PACK instruction...

### Dump Analysis Output Listing

```
0000E0-00001A1F 000008F0 01100700 00000728 00004698 J2L01800 00000028 00284930 *.....0.....-001700
000170-82086404 00280000 000018A8 00000000 00000000 00000000 00270400 *.....-0017F0
000120-00000000 00000000 00000000 00000000 00008668 00000000 00000000 0000185C *.....0.....-001610
000140-00001846 00000000 00000000 0709C6C9 03C54040 80000048 00000000 000004F0 *.....PROFILE .....0.....-001630
000160-F1200000 00000000 00000698 01001920 00000028 002C488F 00884004 00280000 *.....0.....-001850
000180-010014F8 00000000 00000000 00000600 00000000 00000000 00000000 *...8.....-001870
0001A0-00000000 00000000 00000000 00000000 00000000 00000000 F2F3F240 F3F6F540 *.....232 365 -001890
0001C0-000002C5 0303C5E8 68A0E3D6 04404640 40404040 40404040 40404040 40404040 * KELLEY, TOM -001880
0001E0-F1F0F340 40F3F0F2 47A0E5C9 05C3C5D5 E36846C7 C5D6D9C7 C5404040 40404040 *103 362 VINCENT, GEORGE ⑪ -001800
```

These analyses assume that you have a knowledge of machine language, especially main storage addressing. If you need to learn about how machine instructions in general are executed, refer to the appropriate processor programmer reference. To learn what an individual machine instruction does, refer to the *Hardware and Software Programming Quick-Reference Guide* (UP-8868).

## 4.2. BAL Dump Analysis

In this subsection, we will analyze an EOJ dump generated from a BAL program called BALOBJ. We will use some of the TCB fields discussed earlier to help uncover the error causing the dump.

Program BALOBJ, run as part of a job named BALJOB, reads numeric data from an input card, adds the numbers, and outputs them to a printer together with a copy of a character string punched on the input card. The program as written is assembled without any errors, but when we run it, it fails, generating an error code of 20 - a program exception. Anticipating the possibility of such an error, we have put an // OPTION DUMP card in the JCL runstream immediately preceding the EXEC instruction that loads and runs BALOBJ. The program exception causes the operating system to dump the job area containing BALOBJ; because the EOJ dump goes directly to the user's log file, there is no need to assign a printer exclusively to the dump.

### 4.2.1. Materials Used

The materials we will use in the analysis are contained in 4.2.4. They include the following as shown in Figure 4-2:

- The BAL source code for BALOBJ
- The linkage editor allocation map
- The EOJ dump, edited for clarity

### 4.2.2. Outline of EOJ Dump

We will use the step-pointer system outlined in 4.1 to help guide our way around the output listings shown in 4.2.4. An important part of using this system is knowing how the EOJ dump is organized. You should compare the examples and chart presented here with the actual dump shown in 4.2.4.

## Dump Analyses

A portion of a typical EOJ dump looks like this:

JOB-RELATIVE ADDRESSES										ABSOLUTE ADDRESSES
FFFF80	40400305	D2C5C4E3	F0F00200	0000002B	FFFFFF7C0	00000C10	FFFFFF9F8	0CE04000		01B580
FFFFA8	000209E0	00000000	00000000	00000240	00000001	0BE00100	00018600	07000C37		01B5A0
FFFFC8	07420000	00000000	00000000	89F00008	12FF88FD	00084740	F01848FD	000658FD		01B5C0
FFFFE8	F00407FF	9500F028	4780F022	0A5458F0	F02077FF	00FEA45J				01B5E0
<b>PROGRAM REGION</b>										
000000	87F0F020	4004F0F7	82240000	00000000	00000000	00000000	00000000	00000000	00000000	01B600
000020	91401048	4780F030	50001058	58001054	90EC000C	50E00010	41A0F40E	4186F49E		01B620
000040	92001032	91401046	4710F056	92131038	96021032	07F89102	10460718	95201031		01B640

LOCATION  
000000

As you can see, each line of the dump contains 32 bytes of data. The location of that data is given by the two addresses flanking each line. Both addresses always refer to the hexadecimal location of the first byte in the line but differ in their addressing: the right-hand address is the address of the location relative to the entire system, while the left-hand address is relative to the program region. Since the job prologue comes before the program region, its left-hand addresses are given as negative (twos complement) values.

The characteristics of the EOJ dump are summarized in Figure 4-1.

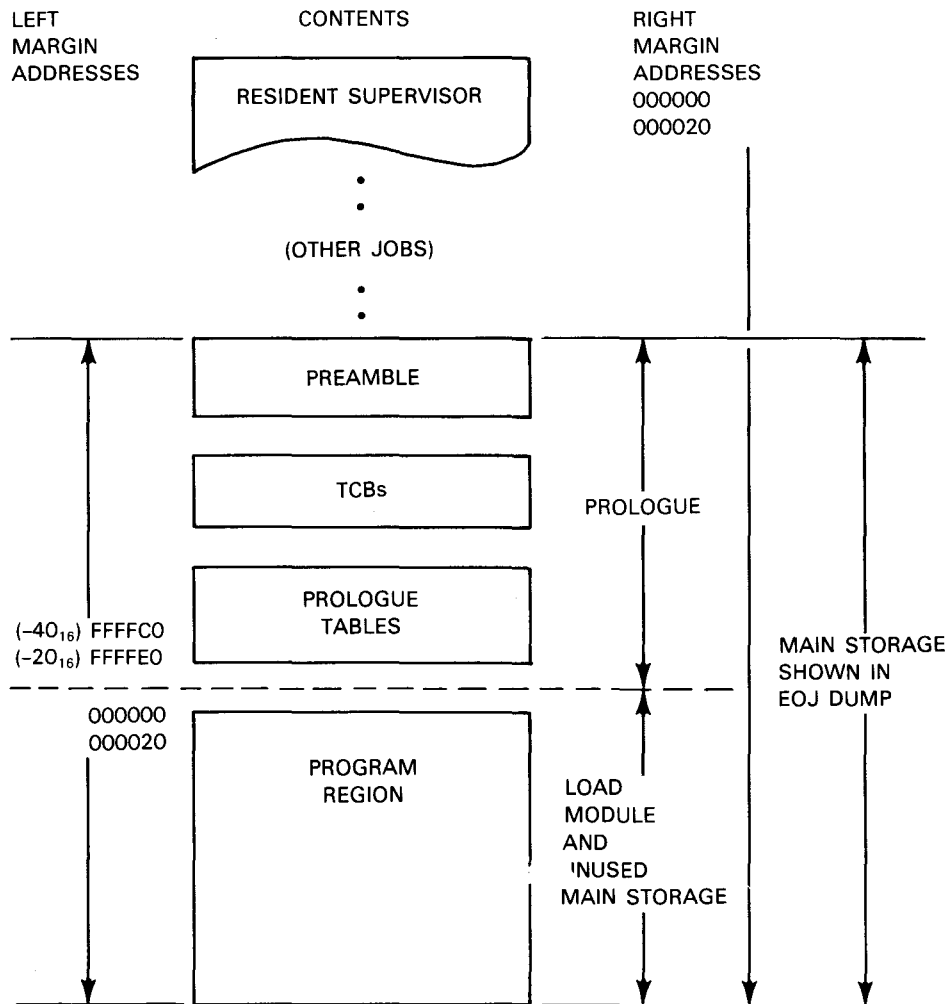


Figure 4-1. EOJ Dump Organization

As Figure 4-1 shows, the EOJ dump presents the entire job region, broken up into a number of blocks: the preamble, each TCB, other prologue tables, and the program region. To read and analyze an EOJ dump, you will need to know about the structure of the TCB, discussed earlier in this manual. As we go through the following analysis you may want to review 2.2 for the structure of the TCB.

## Dump Analyses

Not shown in Figure 4-1 are several other items that are present at the beginning of the EOJ dump:

- Job name
- System version
- PSW at the time the program failed
- Error code
- Address of the TCB that was active when the program failed
- Contents of problem registers 0 through 15 when the program failed

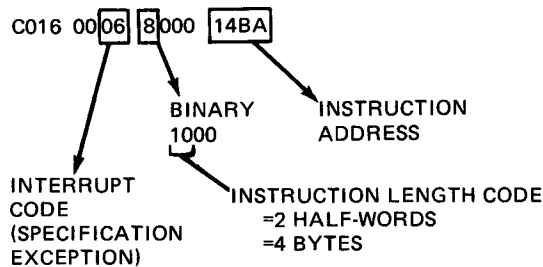
### 4.2.3. Analysis

To determine what caused program BALOBJ to fail, we proceed as follows:

	PSW AT INTERRUPT = <b>C01600068000148A</b>	ERROR CODE = 00000020	TCB ADDR = <b>01AB00</b>
TCBS			
FFF500	1001A000 00000300 2001A000 00000100	0000F22A 0001AA00 1C000000 00000000	<b>01AB00</b>
FFF520	<b>C0160006 8000148A</b> 00000000 80001498	00000000 00000000 00000000 00000000	01AB20
FFF540	00000000 00000000 00000000 00000000	00000000 00000000 00001492 00000000	01AB40
FFF560	00000000 00000000 00000000 00000000	00000000 00000000 00000000 00000000	01AB60
FFF580	00000000 0A540020 6C000000 C0160006	8000148A 00000000 01862800 0C815000	01AB80
FFF5A0	00FED6C8 00000000 00000000 00000000	00000000 00000000 00000000 00000000	01ABA0
FFF5C0	00000000 00000000		01ABC0

We begin by looking for the instruction that directly caused the failure of the program. The TCB at location 1AB00<sub>16</sub> is the primary TCB for the job; therefore, its PSW save area (offset 20<sub>16</sub>-27<sub>16</sub>) ① holds the PSW at the time of the error.

Let's look more closely at the task PSW. Using the PSW illustration in 2.1.1, we can extract the following information:



From this PSW, we can determine where the error occurred and what caused it. We find the error location by remembering that all interrupt handlers store the address of the next instruction to be processed in bits 48-63 of the PSW. The length of the instruction being processed at interrupt time is stored in bits 32-33. From this information, we can find the address of the instruction causing the interrupt:

```

14BA  PSW instruction address
- 4    ILC (= 2 half words)
14B6  Address of instruction causing interrupt
  
```

As to what caused the interrupt, we can determine that from bits 24-31, the interrupt code. In our case, the code is 06, which, according to Table A-1, is a specification exception involving the improper use of the instruction at 14B6.

PROG NAME	TRANS ADDR	FLAG	LABEL	TYPE	ESID	LNK ORG	HIADDR	LENGTH	OBJ ORG
*** END OF AUTO-INCLUDED ELEMENTS ***									
- 79/09/07 21.11 -									
			BALOBJCT	OBJ					
			BALOBJCT	CSECT	01	00001490	0000168F	00000200	00000000
			COFILE	ENTRY	01	00001510			00000380
			COFILEC	ENTRY	01	00001542			00000082
			COFILEE	ENTRY	01	00001548			00000088
			PRFILE	ENTRY	01	00001584			000003F0
			PRFILEC	ENTRY	01	00001582			00000122
			PRFILEE	ENTRY	01	00001588			00000128

To use the information we have so far, we will need to find the instruction that caused BALOBJ to fail. Keep in mind that a load module like the one loaded and run in this job can comprise one or more object modules like BALOBJ, bound together by the linkage editor. We can use the linkage editor allocation map to determine in what portion of the load module our failing instruction lies.

## Dump Analyses

Pointer ② indicates an entry on the allocation map. Three items of the entry are of interest to us: the label BALOBJ, the LNKORG address 1490, and the HIADDR address 168F. The LNKORG and HIADDR addresses define the boundaries of the object module named BALOBJ and we see that address 14B6, indicated by ①, lies within BALOBJ. The exact location of the address can be determined as follows:

```

14B6  Address of failing instruction in load module
- 1490 LNKORG address (first address of BALOBJ)
  26   Offset of failing instruction within BALOBJ
    
```

00001C	9240	C18C	0018C	15	NVI	PRBUFF,C* *	CLEAR OUTPUT BUFFER.	
000020	0226	C180	00180	16	MVC	PRBUFF+1(139),PRBUFF		
000026				17*	GET	CDFILE,CDBUFF	READ INPUT RECORD.	
00002A	5810	C1F0	001F0	18*	DC	CY(0) SET ALIGNMENT		GET00230
00002E	9210	1031	00031	19*	L	1,=A(CDFFILE) LOAD R15, FILENAME ADDRESS		GET00260
				20*	L	C,=A(CDBUFF) LOAD R05, WORKAREA ADDRESS		GET00850
					NVI	49(1),X*10* SET FUNCTION CODE		GET01010

We now know that the failing instruction we seek does lie within our BAL program, rather than within one of the other object modules making up our load module. We turn to the assembler listing for program BALOBJ and look at offset  $26_{16}$  ③, the instruction location calculated in ②. We find there that our failing instruction is one that LOADs register 1 from a storage location determined by base register 12 ( $C_{16}$ ) and offset  $1F0_{16}$ . Looking at Table A-1 for possible causes of our error, we see that the most likely explanation is that the LOAD does not refer to a proper boundary. LOAD instructions can operate only on data residing on full-word boundaries; might the main storage location given by ③ not satisfy this requirement? To see, we go to ④, located at the beginning of the EOJ dump.

PROBLEM	PROGRAM	RESS						
RESS 0-7	00000000	80001498	00000000	00000000	00000000	00000000	00000000	00000000
RESS 8-F	00000000	00000000	00000000	00000000	00001492	00000000	00000000	00000000



To find the storage location referred to by the LOAD instruction in ③, we must first find what is in base register 12 ④. Ignoring the high-order byte, which plays no part in address formation, we see that register 12 contains the value 1492. Once again, the instruction in ③ is:

5810 C1F0

The storage location is determined as follows:

1492 Register 12 contents  
 + 1F0 Displacement  
 1682 Referenced main storage location

Location 1682 does not lie on a full-word boundary; therefore, any LOAD instruction referencing it will generate a program exception like the one that has occurred in this program.

One question remains: why does the main storage location fail to lie on a full-word boundary? As ③ shows, the location is represented by an address constant (for CDFILE) and address constants are assembled on full-word boundaries. For more clues, we will have to look backwards in the program to see what events could have disrupted it.

```

000000 05C0
000000
000002 0700
000004 4510 C010
000008 F0
000009 000080
00000C 80
00000D 0000F0
000010 0A26
000012 9240 C194 00194
000016 0226 C195 C194 00195 00194
00001C 9240 C18C 0018C
000020 0226 C18D C18C 0018D 0018C
000026 5810 C1F0
000026 5810 C1F0

1 BALOBJCT START 0
2 BALR 12,0
3 USING BALOBJCT,12
4 OPEN COFILE,PRFILE
5 CNOP 0,4
6 BAL 1,0(14*3)
7 DC X'F0'
8 DC AL3(CDFILE)
9 DC X'80'
10 DC AL3(PRFILE)
11 SVC 38 ISSUE SVC
12 LOOP MVI CDBUFF,C' '
13 MVC CDBUFF+1(39),CDBUFF
14 MVI PRBUFF,C' '
15 MVC PRBUFF+1(39),PRBUFF
16 GET COFILE,CDBUFF
17 DC (Y(0) SET ALIGNMENT
18 L 1,=(CDFILE) LOAD R15, FILENAME ADDRESS

OPEN FILES.
OPE00130
OPE00500
OPE00580
OPE00600
OPE00550
OPE00600
CLEAR INPUT BUFFER.
CLEAR OUTPUT BUFFER.
READ INPUT RECORD.
GET00230
GET00260
    
```

Looking at ③ again, we see that the displacement 1F0<sub>16</sub> is evenly divisible by 4. For a full-word location, the value contained in register 12 must also be evenly divisible by 4. As ④ shows us, however, the register 12 value of 1492 is not. So we look at previous instructions affecting register 12 and find at ⑤ a BALR instruction and a USING directive.

The USING directive marks the location of symbol BALOBJ as the base address for this assembly and names register 12 as the base register. Since BALOBJ occupies location 00 (1490 when linked into the load module), all references to main storage use register 12, presuming it has a value of 0. However, the previous instruction at ⑤, BALR 12,0 effectively loads a value of 2 (in fact, a load module-relative address of 1492) into register 12. Since the BALR instruction takes effect only at execution time, it introduces a 2-byte address offset that was never taken into account by the assembler. This unexpected offset causes the LOAD instruction in ③ to try to access a main storage location that is not on a full-word boundary. The result is the specification exception that caused BALOBJ to fail.

The solution to the problem is to rearrange or rewrite those introductory BAL statements in ⑤. One correct sequence is:

```
BALOBJ START 0
        BALR 12,0
        USING *,12
        .
        .
        .
```

We should note that the statements at ③ and ⑤ were assembled without error. It was only after they were executed that a problem emerged. We should also note by looking at Figure 4-2 that the steps we take to analyze BALOBJ follow a logical, not a physical, sequence - step ⑤ lies above ③, for example. It is the time sequence that we are most concerned with, and the EOJ dump can be helpful to us in that respect.

### 4.2.4. Dump Analysis Materials

The edited printout from job BALJOB, including the assembler listing for program BALOBJ, the link edit, and the EOJ dump are contained in Figure 4-2. All the pointers referred to in 4.2.3 are shown in the list below in numerical order, each pointer referring to the part of Figure 4-2 on which it appears. Use this list to assist you in looking for these pointers.

- ① part 1
- ② part 7
- ③ part 5
- ④ part 1
- ⑤ part 5



Dump Analyses

FFFFE8	F00407FF	9500F028	4780F022	0A5458F0	F02807FF	00FEA45J	01B5E8
PROGRAM REGION							
000000	47F0F020	40040F07	82240000	00000000	00000000	00000000	01B600
000020	91401048	4780F030	50001058	58001054	90EC000C	50E00010	41A0F40E
000040	92001032	91401046	4710F056	92131038	96021032	07F89102	10460718
000060	4780F206	91801047	4780F07C	9548B1031	4780F120	9548C1031	4786JF11C
000080	4710F08E	921401038	96021032	07FB9588	10314780	F0A2958C	10314780
0000A0	F0849104	104A0788	47F0F04C	9488F104A	95881031	4770F08C	9610104A
0000C0	F0CA5000	102847F0	F0CE9640	104A9101	1C494780	F0E09118	104A4780
0000E0	94F3104A	91181046	4780F0EE	05EA9140	104A4710	F1005820	10285020
000100	928F1J3C	91401047	4710F110	96021047	7A009118	10460776	05EA07F8
000120	948F1047	18229118	10464780	F13005EA	95000016	4780GF146	43200016
000140	0A009640	10479540	00174780	F30E9140	1C474780	F15C05EA	948F1047
000160	4780F168	94F3104A	954C1031	4770F174	96101048	96201047	96201048
000180	4570F18E	0200104E	103C47F0	F3F44130	00089110	1048478J	F1C01923
0001A0	4140F1E6	91071049	4780F1B8	1A239104	1C494780	F1881A23	43224000
0001C0	89200003	41202001	4220103C	91101048	4780F108	96801J3C	91201048
0001E0	9606103C	92021043	94CF1448	07F71109	19210929	39394149	51590969
000200	51590969	39391850	58401050	4690102C	91181046	4780JF21A	05EA4830
000220	10484710	F27E9101	10484710	F2429110	10464710	F25E1864	18694330
000240	F2629548	14444770	F2525830	105847F0	F2624330	104445833	040047FJ
000260	50011839	4780F26A	18334930	104C4700	F27E4830	104C9642	1048964J
000280	10429117	10494770	F29E9101	10434780	F29E4300	10254203	40009A01
0002A0	10474780	F2A454FD	104747F0	F28A9150	10474780	F2C0948F	104702C0
0002C0	91201047	4780F202	96101047	0200103C	104E9110	10464780	F2E81459
0002E0	F2E80630	4430F408	91081046	4780F316	47F0F308	4780F300	9103103C
000300	9106103C	4710F316	5800103C	02021030	10515000	10509104	10484780
000320	91101046	4710F332	06404320	400047F0	F3380650	43205000	417JF396

Figure 4-2. BAL Dump Analysis Listing (Part 2 of 7)

001140	03E09100	10494700	FC664990	FFDA4580	FEB405E8	9120104A	4710FC80	95001059	01C749
001160	4780FC00	4990FF0A	9120104A	4710FC94	9180104B	4710FC94	96031054	91031048	01C760
001180	4780FC00	91021048	4780FC0C	91021049	4780FC0C	45E0F418	45E0F0F0	18631874	01C780
0011A0	47F0FC0C	56201074	45E0F42C	4450F170	18531874	47F0FC0C	91081048	4780FC0C	01C7A0
0011C0	45E0F106	9120104A	4710F0D6	9180104B	4780F066	41201054	5020100C	91041049	01C7C0
0011E0	4780F052	02011061	10560201	105AF064	41401014	50401054	92041054	0A009180	01C7E0
001200	10024710	F0200A01	0201104E	1018947F	104E941F	104F4140	104E5040	10549202	01C800
001220	10589221	10540A00	91801002	4710F04A	0A01D201	10561061	0A749AFF	105E0201	01C820
001240	105A104C	92011054	47F0F0D0	00069102	10464780	F07A9180	10494780	F07A9680	01C840
001260	10035840	10541826	1A469140	10484710	F0B24B40	FDAE0501	4000F080	477CF0CC	01C860
001280	4820FDAE	4920FDAE	4720F08A	47F0F0CC	0C000000	00040002	40404B40	F0ACD503	01C880
0012A0	4000F0A8	4770F0CC	4820FDAC	4920FDAC	4720F0B2	4020105A	0A0047F0	F0F645B0	01C8A0
0012C0	FEA805E8	91801049	4780F0EC	4590FF0A	4590FF58	4580FEC0	94F81049	9180104B	01C8C0
0012E0	4780FE24	0501105E	F44E4770	FE249180	1C024710	FE100AC1	91401002	4710F336	01C8E0
001300	05001026	1027078C	45E0F354	91101046	4710FE4A	9180104B	4780FE3C	45E0F2C8	01C900
001320	47F0FE4A	05E89120	104A4710	FE4A4590	FF0A07FC	96801060	9101104A	078E9110	01C920
001340	10474780	FEA69104	10604710	FE6E9104	10494710	FE869120	10474780	FE7E9610	01C940
001360	106047F0	FEA294EF	106047F0	FEA29120	10474780	FEA29120	10604710	FE9E9620	01C960
001380	106047F0	FEA29610	106094CF	104707FE	4120103C	5020100C	0A0007FB	940F1054	01C980
0013A0	96081054	47F0FE00	94041054	96191054	96081060	47F0FE00	41201054	5020100C	01C9A0
0013C0	91101060	4780FEE8	94EF1060	96201054	91081060	4780FEF8	94F71060	45E0F134	01C9C0
0013E0	0A009120	1060078B	96101060	940F1060	07FB9180	10490789	5840103C	58201078	01C9E0
001400	4450FF52	9120104A	4780FF4C	91041048	4710FF34	41202050	47F0FF38	412020A0	01CA00
001420	9A011066	91031067	47E0FF4C	94FC1067	41201084	50201078	07F90200	20004000	01CA20
001440	91041048	4710FF7A	41401174	41201264	4450FF52	58401054	41201174	4450FF52	01CA40
001460	07F94140	12644120	13044450	FF525840	10544120	12644450	FF5207F9	45E003A8	01CA60
001480	45E00380	45E00380	45E00380	12664770	05C00700	4510C010	F0001510	80001580	01CA80
0014A0	0A269240	C1940226	C195C194	9240C18C	0226C180	C18C5810	C1F05800	C1F49210	01CAA0

Figure 4-2. BAL Dump Analysis Listing (Part 3 of 7)

Dump Analyses

00140	103158F0	103405EF	F222C1E4	C194F222	C1E7C199	FA22C1E7	C1E4F332	C18CC1E7	01C4C0
00140	96F0C18F	D210C1C3	C19E5810	C1F85800	C1FC9220	103158F0	103405EF	47F0C012	01CAE7
00150	5810C1F0	0A275810	C1F80A27	0A1A0700	0C000000	00000000	00000000	00000000	01C800
00150	00000000	00000000	00000000	C3C4C6C9	D3C54040	00000000	00001500	000008F0	01C820
00150	01000000	000004E8	00001500	00001624	0C000000	00700101	82008004	00280000	01C840
00156	00000000	00000000	00000000	00000000	00000000	00280000	0000150C	3C000001	01C860
00156	00000000	00000000	00000000	00000000	0C000000	00000000	00000000	D7D9C6C9	01C880
0015A	03C54040	00000000	00000000	000004F0	F0000000	00700000	00001500	0000164C	01C8A0
0015C	00000000	00000400	C0000004	00280009	00000000	0000150C	00000000	5830001C	01C8C0
0015E	1A300207	30809002	D2076DEE	90020207	79DE9002	921260F6	411060EE	41000001	01CRE0
00160	0A251200	47406622	D2014000	60F20203	4C0860FE	9A014006	920C008E	9218408F	01CC00
00160	47F06640	D2254040	40404040	40404040	40404040	40404040	40404040	40404040	01CC20
00160	40404040	40404040	40404040	40404040	40404040	40404040	40404040	40404040	01CC40
00160	40404040	40404040	40404040	40404040	40404040	40404040	67AC5830	00FCA13D	01CC60
00160	00C01510	00001624	00001580	0000164C	478067AC	953F90C1	47706686	D2076792	01CC80
0016A	900245F0	673A1200	47F06E3E	41106792	4100679A	8900C008	430066CB	41110000	01CCA0
0016C	561066C6	47F066CC	11000000	00230A5E	12004740	67864780	668695C6	67984770	01CCC0
0016E	6E145820	81180507	67A32008	47706772	91802000	47806E9E	48402002	5A468118	01CCF0
00170	43107171	4410678E	47866772	91012000	4710672A	48104004	89100008	5A1L2004	01CD00
00170	41101018	50102004	96012070	58208118	4C203000	41303002	47F066AA	41106792	01CD20
00174	4100679A	89000008	43006758	41110000	56106756	47F0675C	11003000	00630A43	01CD40
00176	12004740	676407FF	41000207	45FC7A5E	47F0601A	41202010	95FF2000	477066E4	01CD60
00178	41000204	47F06E26	41000205	47FC7A56	91004001	00000004	050CF0F0	43900000	01CD80
0017A	00000000	00000016	90C2C103	06C11899	91400078	47806702	948FD078	07017A66	01CDA0
0017C	7A6645F0	7A420501	7A667A68	470087C0	47F06808	91200078	47806804	45F07A42	01CDC0
0017E	95089001	477067F2	95009006	477067DA	47F0685E	12114780	6C449509	90014780	01CDE0
00180	6A6247F0	670A45F0	7A421822	43209001	89200002	44026816	47F06C44	47F06804	01CE00
00180	47F0688E	47F06C44	47F06804	47F06804	47F06804	47F0698C	47F06856	47F06A62	01CE20

Figure 4-2. BAL Dump Analysis Listing (Part 4 of 7)

LOC.	OBJECT CODE	ADDR1	ADDR2	LINE	SOURCE	STATEMENT	
000000				1	BALOBJ	START C	
000000	05C0			2		BALR 12,0	
000000				3		USING BALOBJ,12	
				4		OPEN CDFILE,PRFILE	OPEN FILES.
000002	0700			A 5+		CNOP C,4	OPE0013D
000004	4510	C010		A 6+		BAL 1,+(4*3)	OPE00500
000008	F0			A 7+		DC X'F0'	OPE00540
000009	00008D			A 8+		DC AL3(CDFILE)	OPE00600
00000C	80			A 9+		DC X'80'	OPE00550
00000D	0000FD			A 10+		DC AL3(PRFILE)	OPE00600
000010	0A26			A 11+		SVC 38 ISSUE SVC	OPE00600
000012	9240	C194	00194	12	LOOP	MVI CDBUFF,C'	CLEAR INPUT BUFFER.
000016	D226	C195	C194	00195	00194	MVC CDBUFF+1(39),CDBUFF	
00001C	9240	C18C		14		MVI PRBUFF,C'	CLEAR OUTPUT BUFFER.
000020	D226	C18D	C18C	0018D	0018C	MVC PRBUFF+1(39),PRBUFF	
				16		GET CDFILE,CDBUFF	READ INPUT RECORD.
000026				A 17+		DC (Y(0) SET ALIGNMENT	GET00230
000026	5810	C1F0		001F0	A 18+	L 1,=A(CDFILE) LOAD R15, FILENAME ADDRESS	GET00260
00002A	5800	C1F4		001F4	A 19+	L C,=A(CDBUFF) LOAD R05, WORKAREA ADDRESS	GET00850
00002E	9210	1031		00031	A 20+	MVI 49(1),X*10' SET FUNCTION CODE	GET01010
000032	58F0	1034		00034	A 21+	L 15,524,1) LOAD ADDR OF COMMON I/O	GET01020
000036	05EF			A 22+		BALR 14,15 LINK TO COMMON	GET01030
000038	F222	C1E4	C194	001E4	00194	PACK PADD1(3),ADD1(3)	CONVERT 1ST NUM TO PACKED DEC.
00003E	F222	C1E7	C199	001E7	00199	PACK PADD2(3),ADD2(3)	CONVERT 2ND NUM TO PACKED DEC.
000044	FA22	C1E7	C1E4	001E7	001E4	AP PADD2(3),PADD1(3)	ADD NUMBERS.
00004A	F332	C18C	C1E7	0018C	001E7	UNPK SUM(4),PADD2(3)	CONVERT SUM TO EBCDIC.
000050	96F0	C18F		0018F		OI SUM+3,X'F0'	MAKE LOW-ORDER DIGIT PRINTABLE.
000054	D21D	C1C3	C19E	001C3	0019E	MVC NAMEOUT(30),NAMEIN	MOVE CHARACTER STRING TO OUTPUT
				29		PUT PRFILE,PRBUFF	WRITE OUTPUT RECORD
00005A				A 30+		DC (Y(0) SET ALIGNMENT	GET00230
00005A	5810	C1F8		001F8	A 31+	L 1,=A(PRFILE) LOAD R15, FILENAME ADDRESS	GET00260
00005E	5800	C1FC		001FC	A 32+	L C,=A(PRBUFF) LOAD R05, WORKAREA ADDRESS	GET00850
000062	9220	1031		00031	A 33+	MVI 49(1),X*20' SET FUNCTION CODE	GET01010
000066	58F0	1034		00034	A 34+	L 15,524,1) LOAD ADDR OF COMMON I/O	GET01020
00006A	05EF			A 35+		BALR 14,15 LINK TO COMMON	GET01030
00006C	47F0	C012		00012	A 36+	B LOOP	LOOP BACK.
				37 *			
000070				A 39+ENDPROC		CLOSE CDFILE,PRFILE	END-OF-PROGRAM HOUSEKEEPING.
000070	5810	C1F0		001F0	A 40+	DC (Y(0)	OPE00100
000074	0A27			A 41+		L 1,=A(CDFILE) LOAD R15, FILENAME ADDRESS	OPE00720
000076	5810	C1F8		001F8	A 42+	SVC 39 ISSUE SVC	OPE00730
00007A	0A27			A 43+		L 1,=A(PRFILE) LOAD R15, FILENAME ADDRESS	OPE00720
				A 44+		SVC 39 ISSUE SVC	OPE00730
				A 45+		EOJ	
00007C				A 46+		DS CH	EOJ00850
00007C	0A1A			A 46+		SVC 26	EOJ00870
				47 *			
				48	CDFILE	DTFCD BLKSIZE=40,EOFAADR=ENDPROC,ERROR=ENDPROC,IOAREA1=CDBUFF,X RECSIZE=40,SAVAREA=SAVE	

Figure 4-2. BAL Dump Analysis Listing (Part 5 of 7)

UNIVAC SYSTEM 05/3 LINKAGE EDITOR  
DATE- 79/09/07 TIME- 21.12

VER780403

CONTROL STREAM ENCOUNTERED AND PROCESSED AS FOLLOWS-

```

      /S
      LOADM BALOAD
BALOJECT*RUN LIBE MODULE*
DPSCOM0 *AUTO-INCLUDED*
DRSCOM1 *AUTO-INCLUDED*
    
```

\*DEFINITIONS DICTIONARY\*

SYMBOL.	TYPE.	PHASE.	ADDRESS.	SYMBOL.	TYPE.	PHASE.	ADDRESS.	SYMBOL.	TYPE.	PHASE.	ADDRESS.
BALOJECT	CSECT	ROOT	00001490	CDS10DJ	CSECT	ROOT	000004E8	CDFILE	ENTRY	ROOT	00001510
CDFILEC	ENTRY	ROOT	00001542	CDFILEE	ENTRY	ROOT	00001548	DPSCOM0	ENTRY	ROOT	00000000
DPSCOM1	ENTRY	ROOT	00000000	DPSCOM2	ENTRY	ROOT	00000000	DPSCOM3	ENTRY	ROOT	00000000
DPSCOM4	ENTRY	ROOT	00000000	DPSCOM5	ENTRY	ROOT	00000000	DPSCOM6	ENTRY	ROOT	00000000
DPSCOM7	ENTRY	ROOT	00000000	DRSCOM1	ENTRY	ROOT	000004E8	DRSCOM12	ENTRY	ROOT	000004E8
DRSCOM15	ENTRY	ROOT	000004E8	DRSCOM2	ENTRY	ROOT	000004E8	DRSCOM3	ENTRY	ROOT	000004E8
DRSCOM6	ENTRY	ROOT	000004E8	DRSCOM7	ENTRY	ROOT	000004E8	DRSCOM8	ENTRY	ROOT	000004E8
KESALP	ENTRY	ABS	00001690	KESRES	ENTRY	ABS	00001690	PR\$IOE	CSECT	ROOT	00000000
PRFILE	ENTRY	ROOT	00001580	PRFILEC	ENTRY	ROOT	00001582	PRFILEE	ENTRY	ROOT	00001588

\*\* ALLOCATION MAP \*\*

PHASE NAME	TRANS	ADDR	FLAG	LABEL	TYPE	ESID	LNK ORG	HTADDR	LENGTH	OBJ ORG
BALOAD00		00001690					00000000	0000169F	00001690	
*** START OF AUTO-INCLUDED ELEMENTS -										
- 03/03/78 37.22 -										
				PR\$IOE	OBJ					
				PR\$IOE	CSECT	01	00000000	000004E3	000004E4	00000000
				DPSCOM7	ENTRY	01	00000000			00000000
				DPSCOM0	ENTRY	01	00000000			00000000
				DPSCOM1	ENTRY	01	00000000			00000000
				DPSCOM6	ENTRY	01	00000000			00000000
				DPSCOM2	ENTRY	01	00000000			00000000
				DPSCOM5	ENTRY	01	00000000			00000000
				DPSCOM4	ENTRY	01	00000000			00000000
				DPSCOM3	ENTRY	01	00000000			00000000
- 77/09/29 12.36 -										
				CDS10J	OBJ					
				CDS10DJ	CSECT	01	000004E8	0000148F	00000FA8	00000000
				DRSCOM15	ENTRY	01	000004E8			00000000
				DRSCOM1	ENTRY	01	000004E8			00000000
				DRSCOM2	ENTRY	01	000004E8			00000000
				DRSCOM3	ENTRY	01	000004E8			00000000
				DRSCOM6	ENTRY	01	000004E8			00000000
				DRSCOM7	ENTRY	01	000004E8			00000000
				DRSCOM8	ENTRY	01	000004E8			00000000
				DRSCOM12	ENTRY	01	000004E8			00000000

Figure 4-2. BAL Dump Analysis Listing (Part 6 of 7)



PHASE NAME	TRANS ADDR	FLAG	LABEL	TYPE	ESID	LNK ORG	HIADDR	LENGTH	OBJ ORG
*** END OF AUTO-INCLUDED ELEMENTS -									
- 79/09/07 21.11 -									
			BALOBJ	OBJ					
			BALOBJ	CSECT	01	00001490	0000168F	00000290	00000000
			CDFILE	ENTRY	01	00001510			00000080
			CDFILEC	ENTRY	01	00001542			00000082
			CDFILEE	ENTRY	01	00001548			00000088
			PRFILE	ENTRY	01	00001580			000000F0
			PRFILEC	ENTRY	01	00001582			00000122
			PRFILEE	ENTRY	01	00001588			00000128
00001490									

FLAG CODES -

B - BLM DATA CSECT	D - AUTO-DELETED	E - EXCLUSIVE 'A' REF	G - GENERATED EXTRN	I - INCLUSIVE 'V' REF
L - DEFERRED LENGTH	M - MULTIPLY DEFINED	N - NOT INCLUDED	P - PROMOTED COMMON	R - SHARED REC PRODUCED
S - SHARED ITEM	U - UNDEFINED REF	V - VCON ITEM		

\*ANY OTHER CODES REPRESENT PROCESS ERRORS\*

LINK EDIT OF 'BALOAD' COMPLETED  
 DATE- 79/09/07 TIME- 21.12  
 ERRORS ENCOUNTERED- 0000 UPSI- X'00'

Figure 4-2. BAL Dump Analysis Listing (Part 7 of 7)

### 4.3. COBOL Dump Analysis

In this subsection, we analyze a JOBDUMP generated from a COBOL program called COBOBJ. This program acts much like the BALOBJ program in 4.2: it reads two numbers off a punched card, adds them, and outputs the sum to the printer. In our situation, the program aborts after reading the third data card in the deck. We have inserted an // OPTION JOBDUMP card immediately preceding the EXEC instruction for COBOBJ. We have also, as required, routed the JOBDUMP output to a separate printer file with an LFD of PRNTR, thus causing the JOBDUMP to appear after the entire user log has been output.

#### 4.3.1. Materials Used

The materials we will use in this analysis are contained in 4.3.4. They are as shown in Figure 4-4:

- The COBOL source code for COBOBJ
- The COBOL object map produced by the COBOL compiler in response to the LST=O parameter
- The linkage editor allocation map
- The edited JOBDUMP

#### 4.3.2. Outline of JOBDUMP

As before, we will use the step-pointer method to find our way through our dump analysis. Before we begin analysis, you should look at the portion of the JOBDUMP shown in Figure 4-4. As you can see, these pages interpret the raw dump and present pertinent information in the form of charts, tables, and other narrative.

The remainder of the JOBDUMP shows the contents of the user job region in hexadecimal. Like the EOJ dump shown in 4.2, both margins contain the address of the first byte in each line. Unlike the EOJ dump, though, the JOBDUMP interprets dump data in EBCDIC in a column to the right of the data itself. Also, unlike the EOJ dump, the JOBDUMP uses different base addresses according to the following scheme, shown in Figure 4-3.

As you can see, the simple EOJ dump division of a job into its prologue and its program region is extended in the JOBDUMP. Each table in the prologue has a header and its own self-relative addresses in its left margin. Likewise, each CSECT (control section) within the program region begins with a header and has its own set of self-relative addresses. (Certain spooling and related tables in the prologue keep their negative job-relative addresses in a manner like that of the EOJ dump.)

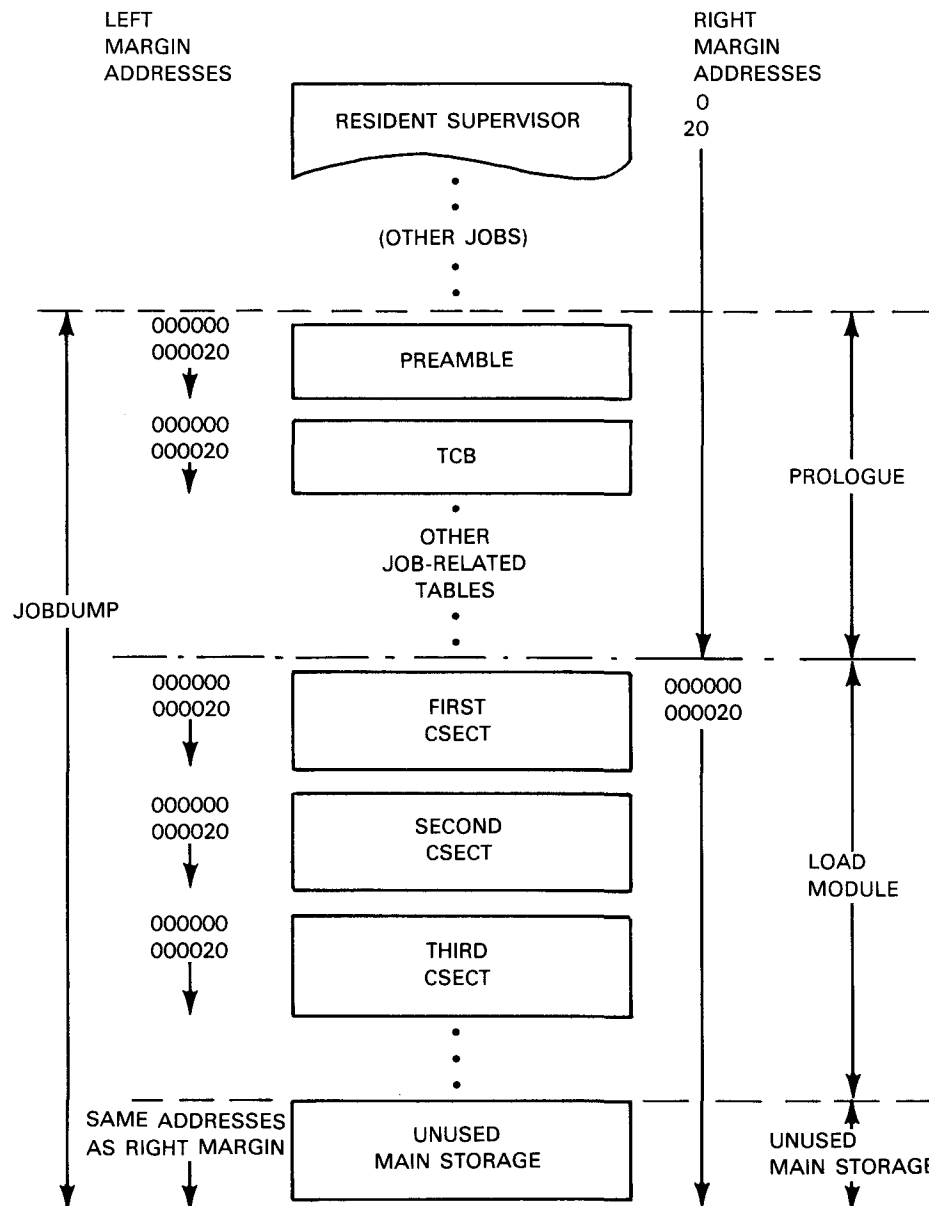


Figure 4-3. JOBDUMP Organization

The prologue/program region division of the EOJ dump is carried over to the right margin except that the JOBDUMP prologue contains absolute system addresses while the JOBDUMP program region has self-relative addresses extending down the entire region. Unused main storage within the job region (that main storage not contained within the program region) is addressed relative to the job in both margins.

The JOBDUMP margin addresses will find much use in the upcoming analysis, so look at the JOBDUMP, Figure 4-4 (parts 6-14), to familiarize yourself with the JOBDUMP addressing scheme.

JOBDUMP requires space on the first VTOC cylinder of SYSRUN to allocate files. Users should eliminate unused files on SYSRUN.

### 4.3.3. Analysis

To determine what caused the JOBDUMP we proceed as follows:

```
-----  
1          TASK CONTROL BLOCK          1  
*          *          *          *          *  
1          1          1          1          1  
-----  
  
TASK CONTROL BLOCK AT ADDRESS 014110  
TASK KEY = 1  
ONLY TASK AT PRIORITY 8  
TCB FLAGS  
  WAIT FOR TRANSIENT  
  WAIT FOR CANCEL IN PROGRESS  
  ORIGINAL PSM HAS BEEN SAVED  
PREAMBLE ADDRESS = D1400C  
TRANSIENT ID/SVC CODE = 1C  
  
① *** TASK PSM ***  
PROGRAM STATUS WORD = C0160007 00001A44  
PROGRAM KEY = 1, WHICH IS JOB COBJOB  
INTERRUPT CODE = 07  
CONDITION CODE = 1  
INSTRUCTION ADDRESS = 001A44  
NONZERO INSTRUCTION LENGTH [16 BYTES]  
OPERATION: [AP] INSTRUCTION: [FA22 AL08 A018]
```

We first look at the task PSW in the JOBDUMP narrative ①. There, all information contained in the PSW is extracted for you. Using this information we see that:

- The interrupt code of 07 indicates a data exception (see Table A-1).
- The failing instruction is add decimal (AP).

- The instruction address can be found by subtracting the AP instruction length from the address given in the PSW:

```

1A44   PSW address
-   6   Instruction length
-----
1A3E   Failing instruction address
    
```

As before, all addresses given are relative to the job base address. Since data exceptions are associated with operands in main storage, we should find the operands addressed by the AP instruction.

REG 0	REG 1	REG 2	REG 3	REG 4	REG 5	REG 6	REG 7
FFFFFFFF	000017A8	0000177C	000005C8	8000055A	020018A8	010018F8	00001948
REG 8	REG 9	REG A	REG B	REG C	REG D	REG E	REG F
007FFFFFFF	007FFFFFFF	00001718	00002718	400019FA	0000199F	40001A10	40001A10

The two operands of the AP instruction shown in ① are both of register-displacement type using register 10 (A<sub>16</sub>). ② points to the contents of register 10 at the time the exception occurred. From this we can find the two operand addresses:

1718	Register 10 contents	1718	Register 10 contents
+ 8	Operand 1 displacement	+ 18	Operand 2 displacement
1720	Operand 1 address	1730	Operand 2 address

```

*** COB0BJVD CSECT COBLODDU PHASE ***
000000-05F045ED F00607FC 98ADF012 9859A060 04A007FE 00001718 00002718 000019E0 *..0.....0.....-0016F0
000020-0000199F 00000000 00000000 00000000 00650436 96F7E59D 927A3002 927A3005 *.....0V.....-001710
000040-00232C00 E598D201 3003E59A 02013096 00001A6E 000019F8 000017A8 000017A8 *...V.K...V.K.....>...8.....-001730
000060-00001820 00001820 00000600 00000508 000004C8 000004E4 007FFFFFFF 00001718 *.....M...U..-001750
    
```

In ③ we see the two AP operands at the addresses calculated in ②. Note that the AP instruction operates on three bytes beginning with each operand. In decimal instructions such as AP, the low-order half byte of each operand is treated as the sign for that operand. Operand 2, at 1730, has a sign value of C, indicating a positive number. Operand 1, however, has an invalid sign value of 4. We can see now that this invalid sign caused a data exception when the AP instruction attempted to operate on it.

# Dump Analyses

At this point, we have found the immediate cause of the error, an erroneous sign. We now attempt to find how the error occurred.

PHASE NAME	TRANS ADDR	FLAG	LABEL	TYPE	ESID	LNK ORG	HIADDR	LENGTH	ORJ ORG
			COBJERR	CSECT	01	0000600	0000727	0000128	0000000
			COBJER1	ENTRY	01	0000674			0000074
			COBJER2	ENTRY	01	0000686			0000086
			COBJER3	ENTRY	01	0000698			0000098
			COBJER4	ENTRY	01	00006AA			000010AA
	- 04/20/79 35.13 -		COBJOJ	OBJ					
			CDS10DJ	CSECT	01	0000728	000016F	00000FC	0000000
			DRSCOM15	ENTRY	01	0000728			0000000
			DRSCOM1	ENTRY	01	0000728			0000000
			DRSCOM2	ENTRY	01	0000728			0000000
			DRSCOM3	ENTRY	01	0000728			0000000
			DRSCOM6	ENTRY	01	0000728			0000000
			DRSCOM7	ENTRY	01	0000728			0000000
			DRSCOM8	ENTRY	01	0000728			0000000
			DRSCOM12	ENTRY	01	0000728			0000000
*** END OF AUTO-INCLUDED ELEMENTS ***	- 79/09/06 00.55 -		COBOBJ00	OBJ					
			COBOBJ00	CSECT	02	000016F	00001A9D	000003AE	0000000
						000016F0			

Because the COBOL compiler generated all the instructions leading up to the error, we now go to the COBOL listing and linkage editor allocation map for program COBOBJ. We will have to look at the allocation map first because the COBOL object program we seek is only one among a number of object modules occupying load module COBLOD. We see from ④ that the AP instruction address of 1A3E lies between the LNKORG address of 16F0 and the HIADDR address of 1A9D, thus placing it inside the COBOBJ00 CSECT. (This CSECT is also labeled in the JOBDUMP; see ③.) We will need to know the displacement of the AP instruction within COBOBJ00 and this can be done as follows:

1A3E Instruction address from ①  
 -16F0 LNKORG address  
 34E AP instruction displacement within COBOBJ00

LINE #	BASE/DISPL	ADDRESS	CONTENTS OF MEMORY	OPERAND ADDRESSES	OPCODE	COMMENTS
00050	C 030	00033A	F2 22 A 018 5 L00	000040	PACK	ADD
	C 036	000340	94 FC A 01A	000042	NI	
⑥	C 03A	000344	F2 22 A 008 5 L05	000030	PACK	
	C 040	00034A	94 FC A 00A	000032	NI	
⑤	C 044	00034E	FA 22 A 008 A 018	000030 000040	AP	
	C 04A	000354	F3 32 6 000 A 008	000030	UNPK	
	C 050	00035A	96 FD 6 003		OI	

Looking at address 34E ⑤ on the COBOL object map, we can see the COBOL instruction that generated the AP. Under the heading OPERAND ADDRESSES are listed the addresses of the two AP operands. We can confirm this by finding their addresses relative to COBOBJ:

16F0	LNKORG address	16F0	LLNKORG address
+ 30	Operand 1 address from ⑤	+ 40	Operand 2 address from ⑤
1720	Operand 1 load module address	1730	Operand 2 load module address

Compare the two addresses to the addresses calculated in ②.

Because operand 1 is the operand that contains the erroneous sign, we should look at instructions prior to the AP instruction that used operand 1 (see illustration above ⑤). Using the OPERAND ADDRESS column, we find that a PACK instruction ⑥ used operand 1 as its destination field. We next look at the source operand.

REG 0	REG 1	REG 2	REG 3	REG 4	REG 5	REG 6	REG 7
FFFFFFF2	000017A8	0000177C	000005C8	8000055A	020018A8	010018F8	000019A8
REG 8	REG 9	REG A	REG B	REG C	REG D	REG E	REG F
007FFFFFFF	007FFFFFFF	00001718	00002718	400019FA	00001990	40001A10	40001A10

⑦

The source field, like the destination field, appears in displacement/base form and uses register 5. Looking at the register contents ⑦, we can calculate the source field address as follows:

(02)0018A8	Register 5 contents (leading byte ignored)
+ 5	Displacement
18AD	Address of source field

For a 4-byte address such as the one in register 5 above, only bits 8-31, the low-order three bytes, are actually used in addressing main storage.

0001A0-00000000	00000000	00000000	00000000	00000000	00000000	F2F3F240	F3F540	.....232	365	-001890	
0001C0-000002C5	03D3C5E8	68A0E3D6	04A0A6A0	40A0A6A0	40A0A6A0	40A0A6A0	40A0A6A0	*	KELLEY, TOM	-001880	
0001E0-F1FDF3A0	40F3F0F2	47A0E5C9	D5C3C5D5	E368A6C7	C5D6D9C7	C5A0A0A0	40A0A6A0	*103	367	VINCENT, GEORGE	-0018D0

⑧

At location 18AD ⑧, we see that the three bytes operated on by the PACK instruction of ⑥ comprise two EBCDIC digits and a blank. The PACK instruction had treated the blank (X'40') as a signed digit by inverting its half bytes (X'40' becoming X'04') thus causing later difficulties by accidentally putting an invalid hexadecimal digit (4) in the sign.

## Dump Analyses

We see by the interpreted data in the right half of the dump printout that these bytes formed part of a character string, very likely part of COBOBJ's data division. To check this we look next at the data division memory map output by the COBOL compiler.

DATA DIVISION MEMORY MAP						
LINE	LEVEL	DATA NAME	REG	DISP	ADDR	LENGTH TYPE
00025	FD	CDFILE				
00028	01	CDBUFF	5	0000		47 GP
00029	02	ADD1	5	0000		3 NUP
00030	02	FILLER	5	0003		2 A/N
00031	02	ADD2	5	0005		3 NUP
00032	02	FILLER	5	0008		2 A/N
00033	02	NAMEIN	5	000A		37 A/N

Looking at the data division main storage map of the COBOBJ compilation listing ⑨, we see that location 18AD lies within the input buffer CDBUFF and has the data-name ADD2. (Compare this address with the source field address in the PACK instruction calculated in ⑦.) We see too that ADD2 is a 3-character unsigned numeric field that is offset within the buffer by five bytes. Let's see what ADD2 contained at the time COBOBJ failed.

ADDRESS	HEX	ASCII	ADDRESS	HEX	ASCII	ADDRESS	HEX	ASCII
0018AD	00000000		00000000	F2F3F290	F3F6F540	.....	232	365 -001890
0001C0	4040D2C5	D3D3C5E8	6840E3D6	04404440	40404440	40404440	40404440	* KELLEY, TOM CDBUFF -001880
0001EC	F1F0F340	40F3F0F2	4740E5C9	D5C3C5D5	E36846C7	C5D6D9C7	C540409C	40404440 *103 302 VINCENT, GEORGE -0018D0

The interpreted data at 18AD ⑩ shows that input numeral 365 incompletely overlaps data area ADD2, the leading digit falling outside of it altogether. By looking at the COBOBJ data division main storage map, it becomes apparent that the program exception was caused when the numeral to be input as ADD2, 365, was punched one column to the left of where it should have been. Repunching the card correctly will resolve this particular problem. In addition, you should include error handling routines in your program to prevent future recurrences.



#### 4.3.4. Dump Analysis Materials

This subsection contains the edited printout from job COBADD, including the COBOL source program, compilation, and link edit for program COBOBJ. In addition, an edited JOBDUMP produced from the execution of JOBDUMP is shown here. All the pointers referred to in 4.3.3 are shown in numerical order, each pointer referring to the part of Figure 4-4 on which it appears.

- |   |         |   |         |
|---|---------|---|---------|
| ① | part 7  | ⑥ | part 3  |
| ② | part 8  | ⑦ | part 8  |
| ③ | part 12 | ⑧ | part 13 |
| ④ | part 5  | ⑨ | part 4  |
| ⑤ | part 3  | ⑩ | part 13 |

COMPILED BY UNIVAC OS/3E COBOL COMPILER VERSION 06.00/09 DATE 79/09/06 TIME 00.54.48

// PARAM IN=COB SRC/INCPUT

// PARAM LST=(S,O,K,L,C)

SOURCE CREATION DATE 79/05/15 TIME 20.30

LINE NO.	SEQ.	SOURCE STATEMENT	IDEN.	PAGE	00001
00001		IDENTIFICATION DIVISION.	COBSR000		
00002		* THIS PROGRAM ACCEPTS TWO 3-DIGIT NUMBERS	COBSR010		
00003		* FROM A CARD RECORD, ADDS THEM, AND OUTPUTS	COBSR020		
00004		* THE SUM, ALONG WITH CERTAIN IDENTIFYING	COBSR030		
00005		* INFORMATION TAKEN FROM THE CARD, TO	COBSR040		
00006		* THE PRINTER.	COBSR050		
00007		PROGRAM-ID. COBOBJ.	COBSR060		
00008		AUTHOR. SYSTEM PUBLICATIONS.	COBSR070		
00009		*	COBSR080		
00010		ENVIRONMENT DIVISION.	COBSR090		
00011		CONFIGURATION SECTION.	COBSR100		
00012		SOURCE-COMPUTER. UNIVAC-9630.	COBSR110		
00013		OBJECT-COMPUTER. UNIVAC-9630.	COBSR120		
00014		INPUT-OUTPUT SECTION.	COBSR130		
00015		FILE-CONTROL.	COBSR140		
00016		SELECT CDFILE ASSIGN TO CARD-READER.	COBSR150		
00017		SELECT PRFILE ASSIGN TO PRINTER.	COBSR160		
00018		*	COBSR170		
00019		DATA DIVISION.	COBSR180		
00020		FILE SECTION.	COBSR190		
00021		* CDFILE IS THE INPUT FILE; CARD INPUT MUST	COBSR200		
00022		* BE IN THE FOLLOWING FORMAT: FIRST 3-DIGIT NUMERAL, 2 BLANKS,	COBSR210		
00023		* SECOND 3-DIGIT NUMERAL, 2 BLANKS, CHARACTER	COBSR220		
00024		* STRING UP TO 30 CHARACTERS LONG.	COBSR230		
00025		FD CDFILE	COBSR240		
00026		RECORD 40 CHARACTERS	COBSR250		
00027		LABEL RECORDS ARE OMITTED.	COBSR260		
00028	01	CDBUFF.	COBSR270		
00029		02 ADD1 PIC 9(3).	COBSR280		
00030		02 FILLER PIC XX.	COBSR290		
00031		02 ADD2 PIC 9(3).	COBSR300		
00032		02 FILLER PIC XX.	COBSR310		
00033		02 NAMEIN PIC X(30).	COBSR320		
00034	FD	PRFILE	COBSR330		
00035		RECORD 40 CHARACTERS	COBSR340		
00036		LABEL RECORDS ARE OMITTED.	COBSR350		
00037	01	PRBUFF.	COBSR360		
00038		02 SUM PIC 9(4).	COBSR370		
00039		02 FILLER PIC X(3).	COBSR380		
00040		02 NAMEOUT PIC X(30).	COBSR390		
00041		02 FILLER PIC X(3).	COBSR400		
00042		PROCEDURE DIVISION.	COBSR410		
00043		BEGINPROG.	COBSR420		
00044		OPEN INPUT CDFILE.	COBSR430		
00045		OPEN OUTPUT PRFILE.	COBSR440		
00046		LOOP.	COBSR450		

Figure 4-4. COBOL Dump Analysis Listings (Part 1 of 14)

LINE NO.	SEQ.	SOURCE STATEMENT	IDFN.	PAGE
00047		READ CDFILE AT END GO TO ENDPROC.	COBSR460	00002
00048		MOVE SPACES TO PRBUFF.	COBSR470	
00049	*	ADD THE NUMBERS, PUT RESULT IN OUTPUT BUFFER.	COBSR480	
00050		ADD ADD1, ADD2 GIVING SUM.	COBSR490	
00051	*	MOVE CHARACTER STRING TO OUTPUT BUFFER.	COBSR500	
00052		MOVE NAMEIN TO NAMEOUT.	COBSR510	
00053	*	WRITE OUTPUT RECORD.	COBSR520	
00054		WRITE PRBUFF.	COBSR530	
00055	*	LOOP BACK.	COBSR540	
00056		GO TO LOOP.	COBSR550	
00057	*	HOUSEKEEPING AT END OF PROGRAM.	COBSR560	
00058		ENDPROC.	COBSR570	
00059		CLOSE CDFILE, PRFILE.	COBSR580	
00060		STOP RUN.	COBSR590	

Figure 4-4. COBOL Dump Analysis Listings (Part 2 of 14)





PHASE NAME	TRANS ADDR	FLAG	LABEL	TYPE	ESID	LNK ORG	HIADDR	LENGTH	ORJ ORG
			CO2BJERR	CSECT	01	00000600	00000727	00000128	00000000
			CO2BJER1	ENTRY	01	00000674			00000074
			CO2BJER2	ENTRY	01	00000686			00000086
			CO2BJER3	ENTRY	01	00000698			00000098
			CO2BJER4	ENTRY	01	000006AA			000000AA
	- 04/20/79 35.13 -		CO510J	OBJ					
			CO510DJ	CSECT	01	00000728	000016EF	00000FC8	00000000
			DRSCOM15	ENTRY	01	00000728			00000000
			DRSCOM1	ENTRY	01	00000728			00000000
			DRSCOM2	ENTRY	01	00000728			00000000
			DRSCOM3	ENTRY	01	00000728			00000000
			DRSCOM6	ENTRY	01	00000728			00000000
			DRSCOM7	ENTRY	01	00000728			00000000
			DRSCOM8	ENTRY	01	00000728			00000000
			DRSCOM12	ENTRY	01	00000728			00000000

\*\*\* END OF AUTO-INCLUDED ELEMENTS -  
 - 79/09/06 00.55 -

000016FD

COB0BJ00	OBJ								
COB0BJ00	CSECT	02	000016FD	00001A90	000003AE	00000000			

4

FLAG CODES -  
 B - BLK DATA CSECT    D - AUTO-DELETED    E - EXCLUSIVE \*A\* REF    G - GENERATED EXTEND    I - INCLUSIVE \*V\* REF  
 L - DEFERRED LENGTH    M - MULTIPLY DEFINED    N - NOT INCLUDED    P - PROMOTED COMMON    R - SHARED REC PRODUCED  
 S - SHARED ITEM    U - UNDEFINED REF    V - VCON ITEM

\*ANY OTHER CODES REPRESENT PROCESS ERRORS\*

LINK EDIT OF \*COBL00\* COMPLETED  
 DATE- 79/09/06 TIME- 00.56  
 ERRORS ENCOUNTERED- 0000 UPSI- X\*00\*

Figure 4-4. COBOL Dump Analysis Listings (Part 5 of 14)

UNIVAC OS/3 JOBDUMP  
 DATE: 79/09/06 TIME: 00:56:35

VER781227

\*\*\* USER ERROR CODE 0020, PROGRAM CHECK

OS/3 VERSION 6.0  
 SUPERVISOR CHARACTERISTIC MASK - 73B2  
 HARDWARE CONFIGURATION MASK - E100

```

CCCCCCC      0000000      BBBB88888      JJJ      0000000      BBBB88888
CCCCCCCCC      000000000      BBBB8888888      JJJ      000000000      BBBB8888888
CCCC  CCCC      0000  0000      BBB  BBB      JJJ      0000  0000      BBB  BBB
CCC   CCC       000   000      BBB   BBB      JJJ      000   000      BBB   BBB
CCC   CCC       000   000      BBB8888888      JJJ      000   000      BBB8888888
CCC   CCC       000   000      BBB8888888      JJJ      000   000      BBB8888888
CCC   CCC       000   000      BBB  BBB      JJJ      000   000      BBB  BBB
CCC   CCC       000   000      BBB  BBB      JJJ      000   000      BBB  BBB
CCCC  CCCC      0000  0000      BBB  BBB      JJJJ  JJJJ      0000  0000      BBB  BBB
CCCCCCCCC      000000000      BBBB8888888      JJJJJJJJJ      000000000      BBBB8888888
CCCCCCC      0000000      BBBB8888888      JJJJJJJ      0000000      BBBB8888888
    
```

```

*-----*
I          I
*      KEY 1      *
I          I
*-----*
    
```

JOB NAME IS COBJOB , JOB NUMBER - 4, STEP NUMBER - 3

ALLOCATION MAP

FROM	TO	LENGTH	CONTENTS
14000	1410F	272	PREAMBLE
14110	141D7	200	TCB
141D8	14223	76	JOB ACCOUNTING TABLE
14224	14353	304	DTF ACTIVE LIST
14354	143D8	136	PHASE LOAD TABLE
145B8	14A07	1104	SPOOLING BUFFERS
14A08	14AEB	228	LOG SPOOL CONTROL TABLE
14DBC	14E9F	228	PRINT SPOOL CONTROL TABLE
15000	16A9F	6816	LOAD MODULE AREA
16AAD	1FFFF	38240	UNUSED MEMORY

LAST PHASE LOADED - COBL0000, PHASE DATE - 79/09/06

REGION DATE - 79/09/06 79/249

\*\*\* JOB CONTROL AREA \*\*\*

JOB STEP OPTIONS

Figure 4-4. COBOL Dump Analysis Listings (Part 6 of 14)

OPTION-DUMP  
 OPTION-JOBDUMP

JOB CONTROL FLAGS  
 JOB ABNORMALLY TERMINATED  
 JOB TERMINATION BUSY  
 ROLL OUT OF JOB INHIBITED  
 PRINT SPOOL FILE GENERATED  
 WTL BUFFERS INITIALIZED  
 ACTIVE PHASE TABLE PRESENT

JOB CONTROL INFORMATION  
 JOB SCHEDULING PRIORITY - LOW  
 JOB CONTROL DIRECTORY DISC ADDRESS - 307/ 2/ 1  
 NO. CYLS FOR ROLLOUT/JOBDUMP - 5  
 JOBDUMP COPY RUNLIB DISC ADDRESS - 307/ 2/ 1

\*\*\* LOADER SEARCH TABLE \*\*\*

SEARCH ORDER	LIBRARY NAME	LIBRARY VSN	PUB ADDRESS	FORMAT 2 LABEL CC/M/R	BEGIN BLOCK	SEARCH BYTF
1	SYSLOD	REL060	0028	154/ 0/14	000001	00
2	SYSRUN	REL060	0028	154/ 1/32	000001	27

①  
 \*-----\*  
 1 \* 1  
 \* TASK CONTROL BLOCK 1 \*  
 1 \* 1  
 \*-----\*

TASK CONTROL BLOCK AT ADDRESS 014110  
 TASK KEY = 1  
 ONLY TASK AT PRIORITY 8  
 TCB FLAGS  
 WAIT FOR TRANSIENT  
 WAIT FOR CANCEL IN PROGRESS  
 ORIGINAL PSW HAS BEEN SAVED  
 PREAMBLE ADDRESS = 014000  
 TRANSIENT ID/SVC CODE = 1C

\*\*\* TASK PSW \*\*\*

PROGRAM STATUS WORD = 00160007 00001A44  
 PROGRAM KEY = 1, WHICH IS JOB COBJOB  
 INTERRUPT CODE = 07  
 CONDITION CODE = 1  
 INSTRUCTION ADDRESS = 001A44  
 NONZERO INSTRUCTION LENGTH (6 BYTES)  
 OPERATION: AP INSTRUCTION: FA22 A008 A018

Figure 4-4. COBOL Dump Analysis Listings (Part 7 of 14)



REG 0 FFFFFFF2 REG 8 007FFFFF	REG 1 000017A8 REG 9 007FFFFF	REG 2 0000177C REG A 00001718	REG 3 000005C8 REG B 00002718	REG 4 8000055A REG C 400019FA	REG 5 020018A8 REG D 00001990	REG 6 010018F8 REG E 40001A10	REG 7 00001948 REG F 40001A10
--	--	--	--	--	--	--	--

\*\*\* TERMINATION INFORMATION \*\*\*

TERMINATION SVC UAS4  
 ERROR STATUS CODE 0020  
 ERROR/PSW ADDRESS = 000700

```

-----*
1          1
*      C D F I L E      *
1          1
-----*
  
```

CARD DTFCO AT ADDRESS 0017A8

CCB: 00 00 80 68 00000000 0000 0000 000017E4 000017CE 0000000000 00 00

MODULE FLAGS = 8000  
 PUB AT 0808 (110) FOR READER 1  
 END-OF-FILE ROUTINE ADDRESS = 001A18  
 RECORD LENGTH ADJUSTMENT = 0  
 FUNCTION CODE = 10  
 ERROR FLAGS = 0000  
 COMMON IOCS MODULE ADDRESS = 000728  
 ERROR MESSAGE CODE = 00  
 USER ERROR ROUTINE ADDRESS = 000698

CCW1:

OP-CODE = 02  
 DATA ADDRESS = 001800  
 FLAGS = 0000  
 BYTE COUNT = 40

IN THE SAVE-AREA:

RECORD SIZE REGISTER DISPLACEMENT = 04  
 IOREG DISPLACEMENT = 28

FLAG BYTE 1 = 49  
 FLAG BYTE 2 = 00  
 FLAG BYTE 3 = 82  
 FLAG BYTE 4 = 08  
 FLAG BYTE 5 = 04

I N P U T F I L E

RECORD FORMAT = FIX UNBLOCKED  
 BLOCK SIZE = 40  
 ALTERNATE DATA ADDRESS = 0018A8  
 RECORD LENGTH = 39  
 EOF-MASK-TABLE DISPLACEMENT = 04

FLAGS -  
 STD MODE  
 IOREG SPECIFIED

Figure 4-4. COBOL Dump Analysis Listings (Part 8 of 14)

```

-----*
1          1
*          *
*   P R F I L E   *
1          1
-----*

PRINTER DTFFR AT ADDRESS 001820

CCB: 00 00 80 68 00000000 0000 0000 0000185C 00001846 0000000000 04 60
MODULE IDENTIFICATION FLAGS = 8000
PUB AT 0048 (FFF) FOR PRINTER 3
RECORD LENGTH ADJUSTMENT = 0
FUNCTION CODE = 20
ERROR FLAGS = 0000
COMMON IOCS MODULE ADDRESS = 00000000
ERROR MESSAGE CODE = 00
CCW1:
  OP-CODE = 01
  DATA ADDRESS = 001920
  FLAGS = 0000
  BYTE COUNT = 40
IN THE SAVE-AREA:
  RECORD-SIZE REGISTER DISPLACEMENT = 00
  IOREG DISPLACEMENT = 2C
FLAG BYTE 1 = 48
FLAG BYTE 2 = 80
FLAG BYTE 3 = 00
FLAG BYTE 4 = 88
FLAG BYTE 5 = 40
RECORD FORMAT = FIX UNBLOCKED
BLOCKSIZE = 40
OP-CODE STORAGE = 00
STD OP-CODE = 01
ALTERNATE DATA ADDRESS = 0018F8
FLAGS -
  CONTROL=YES
  IOREG SPECIFIED
  PRINTOV=SKIP

```

Figure 4-4. COBOL Dump Analysis Listings (Part 9 of 14)

```

- - - - -
1  P R O L O G U E   A R E A
1  - - - - -

```

\*\*\* P R E A M B L E \*\*\*

```

000000-C306C201 D6C2404D 00900000 0000C000 0000B600 50015100 50000354 *C0B408 .....6.....-014000
000020-C306C203 D6C4F0F0 C306C203 D6C4F0F0 7909D600 50C00490 70000888 50000000 *C0A10000C08L0000 .....-014020
000040-00000000 00000000 00000000 00000000 00000000 50000000 50000000 .....0-014040
000060-40404040 0790906C 004F00F9 40F7F9F2 F4F90000 20440908 01F20305 01330102 * ..2. .9 79249.....2.....-014060
000080-00000800 00000000 01330100 00010000 00000000 00000000 00000000 *.....5.....-014080
000100-00000000 00000000 00000000 00000000 00000000 00000000 00000000 *.....-014100
000120-00000000 00000000 00000000 00000000 00000000 00000000 00000000 *.....-014120

```

\*\*\* T C B \*\*\*

```

000020-10014110 000000300 20014110 00000100 00002456 50014700 10000000 00000000 .....-014110
000040-C0160007 D001A44 FFFFFFF2 000017A8 0000177C 00000508 8000055A 020018A8 *.....2.....2.....-014130
000060-010018F8 00001948 0077FFFF 0077FFFF 00001718 00002718 400019FA 00001990 *.....2.....-014150
000080-40001A10 40001A10 00000000 00000000 00000000 00000000 00000000 *.....-014170
000100-00000000 0A540020 00000000 C0160007 D001A44 00000000 01330100 50010000 *.....-014190
000120-00FFA3F4 00000000 00000000 00000000 00000000 00000000 00000000 *.....-014180
000140-00000000 00000000 .....-014100

```

Figure 4-4. COBOL Dump Analysis Listings (Part 10 of 14)

```

FFFEA0-0700C37 07420000 00000000 00000000 80000000 000B4100 0B49000B 49000B59 *.....-014EAD
FFFECD-000B000B 0009000B 49000B49 00000000 00000000 00000B81 00000055 00020000 *.....R.....-014ECD
FFFEED-00000B81 00000B80 00000000 00100005 00000000 00C00000 00000000 01000000 *.....-014EED
FFFF00-00000000 00000000 00000300 00000000 00000000 00880210 00400000 00000000 *.....-014F00
FFFF20-00000004 00009BFC FFFF0004 00000000 000145C8 00000001 00070905 E3094740 *.....H.....PRNTR -014F20
FFFF40-40010000 00009B78 40000000 00210000 0000E2E3 C105C4F1 40400305 02C5C4E3 * .....STAND1 LNKEOT-014F40
FFFF60-F0F00200 00000043 FFFFF588 00000C10 FFFFF5C8 00484000 000209E0 00J00000 *0^.....5.....SH.. -014F60
FFFF80-00000000 00000240 00000001 0C280100 00015J00 07000C37 07420000 00000000 *.....&.....-014F80
FFFA00-00000000 89F00008 12FF88F0 00084740 F01848FD 000658FD F04077FF 95J0F058 *.....0.....0...0...00M...0.-014FAD
FFFFC0-4780F022 0A5458FD F02807FF 00FF11AC 9500F02C 4780F00A 0A54907C 000C98AC *..0...00.....0...0...a....-014FC0
FFFFE0-F024180C 189F8890 00181A99 58F0B114 18FC07FF 10C75810 00000280 J0015000 *0.....0.....0.....G.....L.-014FE0

```

```

*-----*
1 PROBLEM REGISTERS 1
*-----*
1 1
*-----*

```

REG D	REG 1	REG 2	REG 3	REG 4	REG 5	REG 6	REG 7
FFFFFFF2	000017A8	0000177C	00000508	8000055A	020018A8	010018F8	00001948
REG B	REG 9	REG A	REG B	REG C	REG D	REG E	REG F
007FFFFF	007FFFFF	00001718	00002718	400019FA	J0001990	40001A10	40001A10

```

*-----*
1 LOAD MODULE AREA 1
*-----*
1 1
*-----*

```

```

*** PRSIOE CSECT, COBLODDU PHASE ***
000000-4780F020 4004F0F7 75130000 00000000 00000600 00000000 00000000 00000000 *.00. .07.....-000000
000020-91401048 4780F030 50001058 58001054 90EC060C 50E00010 41A0F3EE 4180F47E *. ....C.L.....L.....3...4=-000020

```

Figure 4-4. COBOL Dump Analysis Listings (Part 11 of 14)

```

000020-10560201 105AF084 41401014 50401054 92041054 0A009180 10024710 FD470A01 *..K..J... ..E .....-001448
000040-0201104E 1018947F 104E941F 104F4140 104E5040 10549202 10589221 10540A00 *K..+...".+... .+E .....S.....-001468
000060-91801002 4710FD6A 0A010201 10561061 0A749AFF 105ED201 105A104C 92011054 *.....K...../.....;K...J.<.....-001488
000080-47F0F0F0 00069102 10464780 FD9A9180 10494780 FD9A9680 10035840 10541826 *.P..P.....-0014A8
0000A0-1A469140 10484710 F0024840 FDCE0501 4000FDD0 4770FDEC 4820FDCE 4920FDCE *... ..K. .N. ....-0014C8
0000C0-4720FDAA 47F0FDEC 00000000 00040602 40404640 FDCC0503 4000FDC8 4770FDEC *.....P..... .N. ..H.....-0014E8
0000E0-4820FDCC 4920FDCC 4720FDD2 4020105A 0A0047F0 FE1045B0 FEC805E8 91801049 *.....K ..J...P.....H.Y.....-001508
000100-4780FE0C 4590FF2A 4590FF78 4580FEE0 94FB1049 91801048 4780FE44 0501105E *.....N.;-001528
000120-F4524770 FE449180 10024710 FE300A01 91401002 4710F336 05001026 1027078C *4.....3.N.....-001548
000140-45E0F354 91101046 4710FE6A 91801048 4780FE5C 45E0F2C8 47F0FE6A 05E89120 *..3.....*..2H.P...Y...-001568
000160-104A4710 FE6A4590 FF2A07FC 96801060 9101104A 078E9110 10474780 FEC69104 *..[.....-...[.....F...-001588
000180-10604710 FE8E9104 10494710 FE469120 10474780 FE9E9610 106047F0 FEC294EF *..-.....-..P.B...-0015A8
0001A0-106047F0 FEC29120 10474780 FEC29120 10604710 FE8E9620 106047F0 FEC29610 *..-..P.R.....R.....-..P.R...-0015C8
0001C0-106094CF 104707FE 4120103C 5020100C 0A0007FB 94CF1054 960B1054 47F0FEF0 *..-.....E.....P..7-0015E8
0001E0-94041054 96191054 96081060 47F0FEF0 41201054 5020100C 91101060 4780FF08 *.....-..P.P...E.....-001608
000200-94EF1060 96201054 91081060 4780FF18 94F71060 45ECF134 0A009120 10600768 *.....7..-1.....-001628
000220-96101060 94DF1060 07FB9180 10490789 5840103C 58201078 4450FF72 9120104A *.....-.....E.....[001648
000240-4780FF6C 91041048 4710FF54 41202050 47F0FF58 412020A0 9A011066 91031067 *...Z.....E.P.....-001668
000260-47E0FF6C 94FC1067 41201084 50201078 07F9D200 20004000 91041048 4710FF9A *...Z.....E.....9K... ..-001688
000280-41401174 41201264 4450FF72 58401054 41201174 4450FF72 07F94140 12644120 *.. ..E... ..E...9. ....-0016A8
0002A0-13044450 FF725840 41606004 47F0F5F6 FF7207F9 007892F5 74A34110 749E41F0 *...E... ..-..056...9...S.....P-0016C8
0002C0-E0004190 E00047F0 *.....P .....-0016E8

```

```

*** COBOL JOB SELECT, COBLODDO PHASE ***
000000-05F045E0 F00607FC 98A0F012 98594060 04A007FE 00001718 00002718 000019E0 *.P..D.....D.....-0016F0
000020-00001990 00000000 00000000 00000000 00650466 96F0E590 927A3002 927A3005 *.....0V.....-001710
000040-00232000 E5980201 3003E59A 02013006 00001A6E 000019F8 000017A8 000017A8 *...V.K...V.K.....>...8.....-001730
000060-00001820 00001820 00000600 00000508 000004C8 000004E4 007FFFFF 00001718 *.....M...U...-001750

```

Figure 4-4. COBOL Dump Analysis Listings (Part 12 of 14)



```

-----*
1          1
*      U N U S E D   M E M O R Y      *
1          1
-----*

001A00-103158FD 103405EF 18E2182D 07FED203 1038C34C 5820106D 02032010 C4645820 *...0.....S.....K...C...-K...D...-001A00
001AC0-105CD203 2010C35C 17FF43FD 8U115800 000498EE 000C980C 001492FF 000C07FE *.*K...C*...0.....-001AC0
001AE0-00000000 40000BB8 00000000 40001AA8 00FF11A0 020005E4 40000478 80001980 *... ..U .....-001AE0
001B00-00000016 00000005 00001868 00000000 0000189E 8000195C 000000F7 00001879 *.....*...7...-001B00
001B20-00000000 00001810 00000000 00000000 00000000 00000000 00000000 00000000 *.....-001B20
001B40-00000000 00000000 00000000 00000000 00000000 00000000 00000180 00000FFC *.....-001B40
001B60-00000001 00001893 0000052A 00004A50 00000002 000005E4 40130202 08000000 *.....E...U .....-001B60
001B80-03A20A1A 41000003 58F0A040 05EFOE03 00030300 00000000 00000002 1F001400 *.....0.....-001B80
001BA0-00000000 00000000 00000000 C5D5C403 C9C24440 43970000 00000000 0000001A *.....ENDLIB .....-001BA0
001BC0-9EC3D6C2 03D6C4FD F0790906 00560800 001A9E40 40404040 40404040 40404040 *.CORL0D00..... -001BC0
001BE0-40404040 40404040 40404040 40404040 40404040 40404040 40171C00 80000000 * .....-001BE0
001C00-00000000 00000000 04C4D709 58C9D6C5 4040171C 00800000 0004C800 00000000 *.....DPR510F .....H.....-001C00
001C20-000040C3 C27CD6D7 C3D3F117 1C008000 00000508 00000000 000000F8 C3C27C07 *.. CB00PCL1.....RCB0P-001C20
001C40-D9E3E6E3 171C0080 00000006 00000000 00000001 28C3D67C C2D1C5D9 09404040 *RTMT.....C00BJERR -001C40
001C60-40404040 40404040 40404040 40404040 40404040 00004A50 C5D5C403 C9C24040 * ..EENDLIB -001C60
001C80-0000165C 020006DF 700245E0 060C02D1 903AD7D6 02049030 C1989101 02B74780 *...*K.0.....0.K...P0K...A...K...-001C80
001CA0-C17CD203 06DC7005 45E0D60C 02079068 07000205 06DC7009 410006DC 41109024 *A0K.0.....0.K...P.K.0.....0.....-001CA0
001CC0-45E0D382 47F0C136 18334140 80034130 36019540 40004780 C12C4140 400147F0 *..L..0A.... ..A.. ..0-001CC0
001CE0-C118443D C192D202 9031C19D 45E0D674 45E0D382 947FD7A1 94AFD7A0 47F0D41E *A...A.K...A...0...L...P...P...0M.-001CE0
001D00-9110D7A0 4710C15E 9610D7A0 411000C4 45E0D5EC 07F34110 000645E0 05EC47F0 *..P...A;..P...D..N..3...0..N..0-001D00
001D20-C06A4110 00D545E0 D5EC9640 D7A047F0 C0C69150 07A04780 C0EC9240 902CD21E *.....N..M.. P..0.F..EP..... ..M.-001D20

```

Figure 4-4. COBOL Dump Analysis Listings (Part 14 of 14)

## 4.4. RPG II Dump Analysis

In this subsection, we analyze a JOBDUMP taken from an RPG II program named RPGOBJ. This program reads two numbers from a card like the programs analyzed in 4.2 and 4.3. Unlike them, it divides one number by the other and prints the results along with a character string taken directly from the input card.

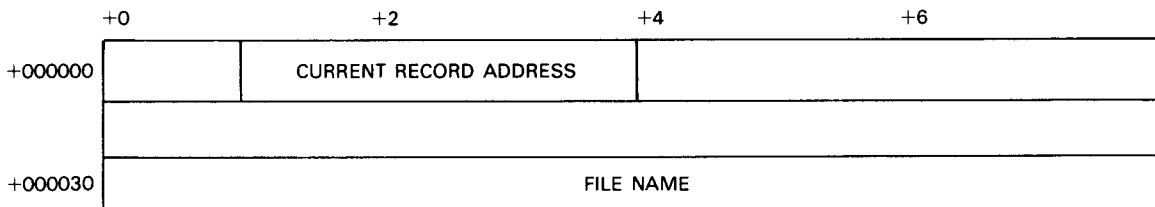
### 4.4.1. Materials Used

The materials used in this analysis are contained in 4.4.5. They are as follows:

- Job log for job RPGCHG (the job which ran RPGOBJ)
- RPG II compilation for RPGOBJ
- Edited JOBDUMP

### 4.4.2. Program-File Interface - the Input/Output Request Block

In this analysis, we will look at a software structure used by RPG II as an interface between an RPG II program and a file. That interface is called the input/output request block (IORB). The IORB is 56 bytes long and contains information that is useful in debugging an RPG II program. For our present purposes the format of the IORB is as follows:



An IORB exists within your RPG II program for each file used by the program. More information on the IORB can be found in the *System Messages Reference Manual* (UP-8076).



### 4.4.3. Program Check Island Code

Another OS/3 facility that we will discuss in this analysis is the program check island code. This code is a software feature that permits users to handle program exceptions with their own routines rather than simply allowing their programs to fail. Several language processors, RPG II included, automatically include island code in load modules, mainly to print error messages and perform other recovery functions when a program fails. As we will see later, the program check island code is designed to save the hardware environment of a failing program in such a way as to let us see what caused the failure. More information on island codes can be found in the *Supervisor Technical Overview* (UP-8831).

### 4.4.4. Analysis

To determine what caused RPGOBJ to fail, we proceed as follows:

```

// OPTION JOBDUMP
// EXEC RPGLOD00
.
.
.
J01 JOB RPGCHG EXECUTING JOB STEP RPGLOD00 0003 21:05:39
RPG030- DIVIDE BY ZERO EXCEPTION
ACT0 LFD - PRTFILE , FORM NAME - STAND1 , COPIES - 0001, PAGES - 00000000, STEP =003
    
```

Here we have loaded and executed our load module RPGLOD00, containing RPGOBJ. The program, however, fails when the RPG030 message ① is printed in the job log. By including an // OPTION JOBDUMP statement in the JCL runstream immediately before the // EXEC RPGLOD00 statement, we can cause OS/3 to generate a JOBDUMP. In doing so we look at the JOBDUMP narrative.

```

-----
1          1
*          *
*          *   TASK CONTROL BLOCK   1   *
*          *
1          1
-----

*** PROGRAM CHECK ISLAND CODE ***

②          BUSY
          ENTRY POINT ADDRESS = 000A70
          PSW/REGISTER SAVE AREA ADDRESS = 000BF0

          PROGRAM STATUS WORD = C016000B E00068C
          PROGRAM KEY = 1 , WHICH IS JOB RPGCHG
          INTERRUPT CODE = 0B
          CONDITION CODE = 2
          INSTRUCTION ADDRESS = 00068C
          NONZERO INSTRUCTION LENGTH (16 BYTES)
          OPERATION: 0P          INSTRUCTION: FDF2 9C17 320F
    
```

At ② we see that a program exception occurred, which activated our program check island code. The major purpose of the island code is to generate the RPG030 message indicated by ①. But the island code also saves the PSW under which the error occurred. We can thus use this PSW to determine that:

## Dump Analyses

- The failing instruction was a divide decimal (DP)
- The interrupt code of 0B indicates an oversize quotient, most commonly generated when dividing by zero

Since operand 2 of the DP instruction represents the denominator, we can confirm that it contains zero, first going to ③.

REG 0 00000008 REG 8 60000672	REG 1 E0000686 REG 9 00000AF0	REG 2 00000AF0 REG A 00000AF0	REG 3 00000000 REG B 000006DA	REG 4 00001000 REG C 0000076C	REG 5 00002000 REG D 00000AF0	REG 6 00001000 REG E 60000044	REG 7 00000000 REG F 41000A72
--	--	--	--	--	--	--	--

Operand 2 of the DP instruction is in base/displacement form using register 3 ③. From register 3, the effective address of operand 2 can then be determined:

0000	Register 3 contents
+ 20F	Displacement
020F	Operand 2 address

Having found its address we now see what operand 2 contains.

```

000200*00000000 00000000 00000000 00890000 00000102 0304*007 090607C5 09E3C9C5 *.....JKLM PROPRIETIE-000200
000220-E2404040 40404040 40404040 40404040 0000000C 40404040 40404040 40404040 *S ..... -000220
    
```

In the JOBDUMP, ④ points to address 20F. The DP instruction specifies that operand 2 is three bytes long, so we highlight addresses 20F-211 to show that operand 2 does indeed contain a packed decimal value of zero.

The next question is how the zero value got into operand 2. To help answer that question, we turn to the RPG II compilation listing that generated RPGOBJ.

FIELD NAMES	ADDRESS FIELD	ADDRESS FIELD	ADDRESS FIELD	ADDRESS FIELD	ADDRESS FIELD
000180 *ERROR	0001A4 PPSDMP EXTRN	00020C ADD1	00020F ADD2	000212 NAMEIN	
000230 QUOTE	000234 NAMEOT				

The field name list in our RPG II listing shows that our operand 2 address of 20F corresponds to a field within our program named ADD2 ⑤. We turn to the RPG II source listing itself to see what ADD2 is.

```

001      000 H D      I
002      001 FCDFILE IPEAF 45 45      READER
003      002 FPRFILE 0 V 48 48      PRINTER
004      003 ICDFILE 011 01 01NCI

004      004 I      1 50ADD1
005      005 I      8 120ADD2
006      006 I      15 44 NAMEIN
    
```

RPG60BJ  
RPG5R010  
RPG5R020  
RPG5R030  
NOTE 117  
RPG5R040  
RPG5R050

⑥

Line 005 of the source program ⑥ shows that ADD2 is a 5-byte decimal number contained in bytes 8-12 of the input record for card file CDFILE (defined in line 003). Knowing this, we can then look for the input card that presented a zero value in the ADD2 field, thus causing the program to fail.

To pick out the error card we first find the data read into the job from that card. The RPG II compiler listing and the *System Messages Reference Manual* (UP-8076) can both help us here. Our intention is to find the IORB associated with file CDFILE. To do so we first look at the RPG II compiler listing under the heading "PROGRAM POINTERS":

```

PROGRAM POINTERS
TABLE INPUT/OUTPUT      CC0254
INPUT FIELD EXTRACTION  CC0258
DETERMINE RECORD TYPE   CC030C
GET INPLT RECORD        CC050C
DETAIL CALCULATIONS     CC067C
TOTAL CALCULATIONS     CC06A2
OVERFLOW OUTPUT         CC06F6
OVERFLOW BYPASS        CC06DA
HEADER/DETAIL OUTPUT    CC0704
OUTPUT FIELDS           CC076C
INPUT/OUTPUT REQUEST BLOCKS  CC08C8
    
```

⑦

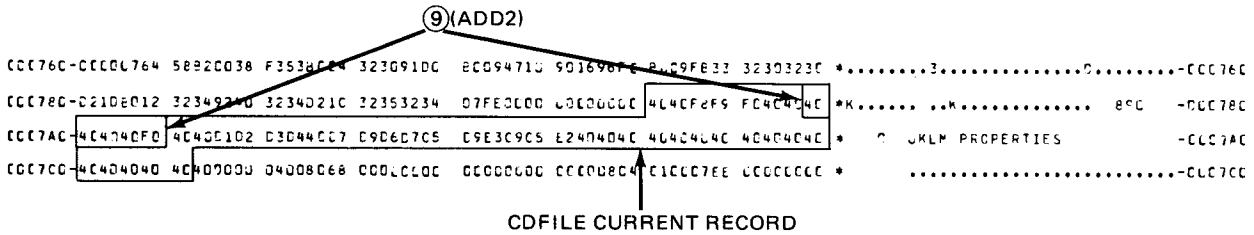
As we can see the input/output request blocks ⑦ are located starting at address 8C8. We go next to that address and follow the IORB format shown in 4.4.2. RPG II always arranges IORBs in the same sequence as the file description statements for their respective files. Because file CDFILE is the first file defined in the program, its IORB occupies the first 56-byte IORB slot, at address 8C8:

```

⑧ IORB CURRENT RECORD ADDRESS      IORB FILENAME
CC08C8-CC0C00400 E7C9C6C9 03C54040 8090151J 0000JF8C 00C104FC F620CCCC 12FFFF0C *...PRFILE .....0.....-CC08E8C
CC0E40-CC0E0F96 01C00832 0E000E3D 0020440C 849A2004 003000J5 00000832 00000000 *.....-CC0E8A0
CC0E80-CC0E0000 42300000 32000758 00000000 00200001 00C100FF FFFFFFF0 00000011 *.....-CC0E8C0
CC0E80-CC0E0000 00000000 00000000 00000000 00000000 00000000 00000000 00000000 *.....-CC0E8E0
    
```

Bytes 1-3 ⑧ of the CDFILE IORB hold address 798, the current record address.

At address 798 we find:



The large block indicates the CDFILE record most recently released to the program, and ⑨ points to bytes 8-12 within the record, the field defined by RPG II field AD2 (see ⑥). We see that ADD2 does contain a value of zero, perhaps a mispunch. We can also tell, from the interpreted dump in the right column, just which input card (JKLM PROPERTIES) contained the faulty data.

#### 4.4.5. Dump Analysis Materials

On the following pages is Figure 4-5, the edited printout from which the dump analysis of 4.4.4 is taken. The following list indicates where in Figure 4-5 each of the pointers of 4.4.4 can be found.

- ① part 1
- ② part 5
- ③ part 6
- ④ part 7
- ⑤ part 2
- ⑥ part 2
- ⑦ part 2
- ⑧ part 9
- ⑨ part 8



```

GCC H 1
001 C10 FCOFILE IFEAF 45 45 READER
002 C2C FPRFILE 0 F 48 48 PRINTER
003 C3C ICCFILE 011 01 CINCI
004 C4C I 1 SCACC1
005 C5C I 8 12CACC2
006 C6C I 15 44 NAMEIN
007 C7C C ADD1 DIV ACC2 QUOTE 6C
008 C8C C MCVE NAMEIN NAMECT 3C
009 C9C CPPFILE 0 1 01
010 1CC C CLCTE B 10
011 11C C NAMECT E 48

```

RPG08J

SYMECL TABLES

RESULTING INDICATORS

ADDRESS RI	ADDRESS RI	ADDRESS RI	ADDRESS RI	ADDRESS RI	ADDRESS RI	ADDRESS RI
CCCC14 1P	CCCC15 LR	CCCC16 00	CCCC17 01	CCCC7A L0	CCCC85 M1	CCCC86 M1
CCCC07 F2	CCCC88 F3	CCCC89 H4	CCCC8A F5	CCCC8F F6	CCCC8C F7	CCCC8D H8
CCCC8E F9	CCCC8F L1	CCCC90 U2	CCCC91 U3	CCCC92 U4	CCCC93 U5	CCCC94 U6
CCCC95 L7	CCCC96 L8					

FIELD NAMES

ADDRESS FIELD	ADDRESS FIELD	ADDRESS FIELD	ADDRESS FIELD	ADDRESS FIELD
CCC180 *ERROR	CCC200 ADD1	CCC20F ACC2	CCC212 NAMEIN	CCC230 QUOTE
CCC234 NAMECT				

002

NOTE 201

NOTE 201 RESULTING INDICATOR IS INVALID OR UNDEFINED. ENTRY OF L0 IS ASSUMED.

PFCCRAM POINTERS

TABLE INPUT/OUTPUT	CCC254
INPUT FIELD EXTRACTION	CCC258
DETERMINE RECORD TYPE	CCC30L
GET INPLT RECORD	CCC500
DETAIL CALCULATIONS	CCC67C
TOTAL CALCULATIONS	CCC6A2
OVERFLOW OUTPUT	CCC6B6
OVERFLOW BYPASS	CCC6DA
HEADER/DETAIL OUTPUT	CCC704
OUTPUT FIELDS	CCC76C
INPUT/OUTPUT REQUEST BLOCKS	CCC8C8

Figure 4-5. RPG Dump Analysis Listing (Part 2 of 9)

URIVAC 05/3 JOBDUMP  
DATE: 80/06/26 TIME: 20:50:33

VER3006C4

\* \* \* USER ERROR CODE 0000

OS/3 VERSION 7.0-C.01  
SUPERVISOR CHARACTERISTIC MASK - 37BFE8AE  
HARDWARE CONFIGURATION MASK - E104

RRRRRRR	FFFFFFF	GGGGGGG	CCCCCCC	HHH	HHH	GGGGGGG
RRRRRRRRR	FFFFFFFFF	GGGGGGGGG	CCCCCCCCC	HHH	HHH	GGGGGGGGG
RRR RRRR	PPP PPP	GGG GGG	CCC CCC	HHH	HHH	GGG GGG
RRR RRR	PPP PPP	GGG	CCC	HHH	HHH	GGG
RRR RRRR	PPP PPPP	GGG	CCC	HHHHHHHHH		GGG
RRRRRRRRR	FFFFFFFFF	GGG	CCC	HHHHHHHHH		GGG
RRRRRRR	FFFFFFFFF	GGG GGGG	CCC	HHH	HHH	GGG GGGG
RRR RRR	PPP	GGG GGGG	CCC	HHH	HHH	GGG GGGG
RRR RRR	PPP	GGG GG	CCC CCC	HHH	HHH	GGG GG
RRR RRR	PPP	GGGGGGGGG	CCCCCCCCC	HHH	HHH	GGGGGGGGG
RRR RRR	PPP	GGGGGGG	CCCCCCC	HHH	HHH	GGGGGGG

```

*--*--*--*--*--*--*--*--*
1                                1
*      K E Y  1                *
1                                1
*--*--*--*--*--*--*--*--*

```

JOB NAME IS R06CH0 , JOB NUMBER - 1, STEP NUMBER - 3

ALLOCATION MAP

FROM	TO	LENGTH	CONTENTS
19A00	19B1F	288	PREAMBLE
19B20	19C47	296	TCB
19C48	19C93	76	JOB ACCCLNTING TABLE
19C94	19C73	224	OPEN FILE TAELE
19C84	19ECB	136	PHASE LOAD TABLE
1A0E4	1A5C3	1104	SPOOLING BLFFERS
1A5C4	1A5E7	228	LOG SPOOL CONTROL TAELE
1A6E8	1A7BB	228	READER SPOCL CONTROL TABLE
1A7EC	1A89F	228	PRINT SPCOL CONTROL TABLE
1AACC	1BAA7	4264	LOAD MODLLE AREA
1BAA8	1FFFF	17752	UNUSED MEMORY

LAST PHASE LOADED - R06L000, PHASE DATE - 80/06/26

REGION DATE - 80/06/26 80/178

\* \* \* J C B C O N T R O L A R E A \* \* \*

JCB STEP OPTIONS

Figure 4-5. RPG Dump Analysis Listing (Part 3 of 9)





```

INSTRUCTION ADDRESS = 00CB18
NONZERO INSTRUCTION LENGTH (2 BYTES)
OPERATION: SVC RPCM2 INSTRUCTION: 0F3B

REG 0      REG 1      REG 2      REG 3      REG 4      REG 5      REG 6      REG 7
0000000B  EC000686  00000B3C  00000000  00001000  00002000  00003000  00000000
REG 8      REG 9      REG A      REG B      REG C      REG D      REG E      REG F
60000672  00000B3C  00000B3C  000006EA  0000076C  00000B3C  60000000  41000A6A

*** PROGRAM CHECK ISLAND CODE ***

BUSY
ENTRY POINT ADDRESS = 00A6E
PSW/REGISTER SAVE AREA ADDRESS = 000030

PROGRAM STATUS WORD = 0016000B EC00068C
PROGRAM KEY = 1 , WHICH IS JCB RPCM2MP
INTERRUPT CODE = 0B
CONDITION CODE = 2
INSTRUCTION ADDRESS = 00068C
NONZERO INSTRUCTION LENGTH (6 BYTES)
OPERATION: 0F INSTRUCTION: FCF2 9C17 320F

REG 0      REG 1      REG 2      REG 3      REG 4      REG 5      REG 6      REG 7
0000080E  800004A8  02000798  00000000  00001000  00002000  00003000  00000000
REG 8      REG 9      REG A      REG B      REG C      REG D      REG E      REG F
60000672  00000B3C  00000B3C  000006EA  0000076C  00000B3C  60000000  00000670

*** TERMINATION INFORMATION ***

TERMINATION SVC 99JF
ERROR STATUS CODE J000
ERRCR/PSW ADDRESS = 019268

*** TERMINATION PSW ***

PROGRAM STATUS WORD = 0016000B EC00068C
PROGRAM KEY = 1 , WHICH IS JCB RPCM2MP
INTERRUPT CODE = 0B
CONDITION CODE = 2
INSTRUCTION ADDRESS = 00068C
NONZERO INSTRUCTION LENGTH (6 BYTES)
OPERATION: 0F INSTRUCTION: FCF2 9C17 320F

*-----*
1 1
* C D F I L E *
1 1
*-----*

CARD DTFCED AT ADDRESS 000708

```

Figure 4-5. RPG Dump Analysis Listing (Part 5 of 9)

```

FFFFAD-0CC00000 89F00008 12FF88F0 00084740 F01848F0 00C658FC FGD4C7FF 58F0FC28 *.....C.....C... 0..C...00P...C0.-01A9AC
FFFFC0-07FFCCCC CCCCC000 0C000CCC 0CFF39CC 9500F02C 47E0F00A 0A540AE9 2C190A19 *.....0...0...Z....-C1A9CC
FFFFE0-0A1C07FF CA270A1C 0C000CCC 0G000000 0C000000 FFFFF374 00000000 0C000000 *.....3.....1....-01A9EC

```

```

*-----*
1                                     1
*   P R O B L E M   R E G I S T E R S   *
1                                     1
*-----*

```

3

REG 0 0C00000E REG 8 6C000672	REG 1 EC0C0686 REG 9 0C0C0B3C	REG 2 0C0C0B3C REG A 0C0C0B3C	REG 3 <u>0D000A00</u> REG B 0D0006EA	REG 4 0G0C10CC REG C 0C0C076C	REG 5 0D0C20CC REG D 0C0C0B3C	REG 6 0C0C3LCC REG E 6CCCC0CC	REG 7 0D0000CC REG F 41000A6A
--	--	--	---	--	--	--	--

```

*-----*
1                                     1
*   L C A D   M O D U L E   A R E A   *
1                                     1
*-----*

```

```

*** R P G C B J      C S E C T ,   R P G L O D U G   P H A S E ***
0C0000-05F058F0 F0C0C7FF 0C000E30 00000L30 0C000000 0C000000 0C000000 0C000000 *..C..00.....00.....-0C0000
0C0020 TC 0C005F   SAME AS LAST WORD

0C0060*0C000000 0C000000 0C000000 0C000000 0C000000 0C000000 0C000000 0C000000 *.....C.....-0C0060
0C0080-0C000000 0C000000 0C000000 0C000000 0C000000 0C000000 0C000000 0C000000 *.....C.....-0C0080
0C00A0-0C000000 0C000254 0C000000 0C000000 0C000000 0C000000 0C000258 0C00030C *.....-0C00A0
0C00C0-0C000000 0C000000 0C000000 0C000000 0C000000 0C00067C 0C0006A2 0C000000 *.....-0C00C0
0C00E0-0C000686 0C0006DA 0C000000 0C000704 0C000000 0C000000 0C000000 0C00076C *.....-0C00E0
0C0100-0C000000 0C000000 0C0008C8 0C000938 0C000938 0C000000 0C000000 0C000000 *.....-0C0100
0C0120 TC 00013F   SAME AS LAST WORD

```

Figure 4-5. RPG Dump Analysis Listing (Part 6 of 9)



```

CC04E0-B0060000 C394D203 1C04401A 45FE0032 B004016C 41A3006E 50A10008 47F0E006 *.....K... ..E.....C..-CC04E0
CC0500-058047F0 80148010 00000000 00000388 CC000000 GC3E181C 580C315C 5010000E *...C.....E.....-000500
CC0520-50E00004 C7C8307B 3C7847F0 80BA58CC 31084AC1 00C2582C 00CC5C2C 317C589E *E...P...#...C.....CA.....E...-CC0520
CC0540-80E61A93 58A10008 96FC0000 41F03085 19AF4770 80E500C 31B892C1 31B091C1 *.....C...C.....>E.....A.....-CC0540
CC0560-318947E0 806E9110 316847EC 806E0A3E 58AC30B8 4A708012 5C1AC000 58F1000C *.....>.....>.....C.....E.....1..-CC0560
CC0580-12FF4780 80921AF3 5EC03000 58AC0184 1AA305EF 58ED0C04 58000006 07FE001E *.....3.....-CC0580
CC05A0-000004A4 00000000 00000000 F0F0F0F0 F0F0F0F0 F0F0F0F0 00000000 5810809E *.....CC00000000.....-CC05A0
CC05C0-1A1358C0 31C84A01 JC061890 58F0800A 1AF305EF 92C0900E 95AA101C 477060EA *.....C.....0...3.....-CC05C0
CC05E0-58F0800E 1AF305EF 47F061C2 58FC311C 05EF9110 901647EC 81C258FC 90C412FF *..C...3...C...C.....C.....-CC05E0
CC0600-47706106 95F03085 58A030B8 4770812C 58A030B8 4A708012 16EE5CEA 000047FC *.....C.....C.....E.....C...-CC0600
CC0620-80929522 9C134770 81564111 000492F0 100158F0 80A212FF 478C8146 50F0809E *.....C...C.....E.....C...-CC0620
CC0640-070367A2 8CA247F0 80BA96FC 30150208 307E8CAA 47F08092 508C0014 58F030B8 *P.....C...C...K...#...C...E.....C...-CC0640
CC0660-41FF0000 05EF5880 JC1447F0 80200000 05805890 3150900E 91640732 90C09000 *.....C.....E.....P.....-CC0660
CC0680-FF029017 320CF0F2 9017320F F63C3230 9017940F 3230021C 32343212 980E9164 *80.....3.....E.....K.....-000680
CC06AC-07F00580 58903150 900E9164 980E9164 419000FF 07FE58AC 315C900E AC00582C *.....E.....E.....-CC06AC
CC06C0-31085800 310C5800 30FC5880 30E40607 30003000 07E8966C 318E947F 318B0580 *.....C.....L0.....".....-CC06C0
CC06E0-95F02043 4770801A 92272047 92022046 41020036 05E0920C 20489200 204090FE *..C.....<...-0006E0
CC0700-ACC007FE 58A03150 900EAC00 58203108 58003100 580030FC 058095FC 3017477F *.....E.....C.....-000700
CC0720-802E9251 20489210 204A5892 00389240 9000022E 90019000 589C0000 05E99202 *.....C.....K.....2...-CC0720
CC0740-20464102 003805E0 058047F0 800E96FC 8001980E ACC007FE 980EAC00 07FE0000 *.....C...C.....-000740
CC0760-00000764 58920038 F3538004 32309100 80094710 901696FC 8009FB33 32303230 *.....3.....0.....-CC0760
CC0780-C210E012 32349240 32340210 32353234 07FE0000 00000000 4040FBF9 FC404040 *K.....K.....800 -000780
CC07AC-404040F0 40405102 03044007 09060705 09E30905 E2404040 40404040 40404040 * .. JKLM PROPERTIES -CC07AC
CC07CC-40404040 40400000 04008068 00000000 00000000 00000804 010007EE 00000000 * .....-CC07CC
CC07EC-00000400 03040609 03054040 80001650 0000135A 000008FC 01100000 19FFFF00 *....CDFILE ...E...I...C.....-CC07EC
CC0800-00001062 02000798 0000002E 00204100 82080204 00200000 00000798 00000000 *.....-CC0800
CC0820-00000000 00000000 00000000 00200400 00104040 4040F0FC F0F0F2F6 40404040 *.....000026 -CC0820
CC0840-40404040 01020340 E309E4C3 02090507 40030648 40404040 40404040 40404040 * ABC TRUCKING CO. -000840
CC0860-40400000 00000000 00000068 00000000 00000000 000008A4 0200088E 00000000 * .....-CC0860

```

Figure 4-5. RPG Dump Analysis Listing (Part 8 of 9)

CCCC880-0CCCC00400	E7C9C6C9	D3C54040	8000151J	00000F8C	00C104FC	F62CCCCC	12FFFFDC	*....PRFILE	.....06.....	-00088C
CCCC8A0-CCCC0CF96	C10C0832	0C0CJC30	002C440C	849A2004	00300C95	00C00832	0C0C0C0C	.....	.....	-0008AC
CCCC8C0-0CCCC0000	423C0000	02000798	00C00000	0C203001	00C100FF	FFFFFFC0	0C00CC11	.....	.....	-0008CC
CCCC8E0-CCCC0000	0C0007C8	00300FAC	0C000000	0C3C4C6C9	03C5404C	.....	.....	IORB FILENAME	CCFILE	-0008EC
CCCC900-0CCCC0832	0C000000	0C300001	000202FF	FFFFFF00	0C000011	1C200000	0C0C0C0C	.....	.....	-00090C
CCCC920-0CCCC0868	0C00CE48	0C000000	0C000000	07C9C6C9	03C5404C	05FC9C42	FC6A187F	.....PRFILE	..C..C..	-00092C
CCCC940-1E409101	40194710	708E9180	401647E0	70269182	40CE478C	7C2E58FC	3CB4C5EF	.....	.....	-00094C
CCCC960-95204018	4770706A	950C40CE	4770706A	58904000	950C4008	4770705C	1B55435C	.....	.....	-00096C
CCCC980-40090650	42507053	92409000	02009001	900047F0	706A1814	58FC3114	05EF057C	.....	.....	-00098C
CCCC9A0-5E707060	58003150	41000110	410070EA	50000004	58104020	5870C024	05E7057C	.....	.....	-0009AC
CCCC9C0-5E707090	47F07090	18FF12FF	478070CA	92F03085	91104016	471070B2	5CF070E6	.....	.....	-0009CC
CCCC9E0-02013181	70E60200	318370E9	984270EA	91013189	47E7F000	0A3B58EC	318C17FE	*K.....WK.....Z.....	.....	-0009EC
CCCCA00-47FC70EC	918C4016	47EC70EC	950C400E	477070EC	58F03084	05EF9842	70EA07FE	.....	.....	-000ADC
CCCCA20-0CCCC0000	0C000100	0C000000	0C000000	0C000000	50C00502	0C000000	0C000000	.....	.....	-000A2C
CCCCA40-0CCCC0832	0C000760	0C000000	50C005F2	50C0093A	0C000000	0C000000	0C000000	.....	.....	-000A4C
CCCCA60-0CCCC0001	0C000000	05F05820	315092F0	30855610	21C45010	31804610	FCBE9505	.....	.....	-000A6C
CCCCA80-21034780	FC320201	318610C2	95072103	4780F052	950B2103	4780F088	41102100	.....	.....	-000A8C
CCCCAA0-5C103188	92F03180	91033189	4770FCAC	9108318A	4710FC9C	47FCFCAC	910C1000	*E.....C.....C.....D.....C.....	.....	-000AAC
CCCCACC-47E0F032	92E73180	91033189	4770F08E	9108318A	4710FC9C	50102104	980F210E	.....	.....	-000ACC
CCCCAEC-0A3BC5EC	5810E008	47F0E000	8C000000	0A0F92E9	31E09103	31894770	FCAC9108	.....	.....	-000AEC
CCCCB00-318A47E0	FCAC5810	31A45010	21045604	318A47F0	FCAC5A3E	5810F0B6	47F0F0BA	.....	.....	-000B0C
CCCCB20-8CCCC0000	0ACF40B8	0C000000	0C000000	0C000000	0C000000	0C000000	0C000000	.....	.....	-000B2C
CCCCB40-0CCCC0000	0C000000	0C000000	0C000000	0C000000	0C000000	0C000000	0C000000	.....	.....	-000B4C
CCCCB60-0CCCC0000	0C000000	0A115810	31504100	0C014120	0C0841EC	308E41F0	0CF04200	.....	.....	-000B6C
CCCCB80-0C539180	10004780	005E42F2	E0008900	0C014620	004E47FC	0070400E	06094001	.....	.....	-000B8C
CCCCBA0-03034006	E4E307E4	E34003C9	18000700	41101000	56100080	47FC0090	40000000	*LL OUTPUT LI.....	.....	-000BAC
CCCCBC0-0A111200	4780009A	0A100504	1000019A	477000AE	56113185	02010100	02800505	.....	.....	-000BCC
CCCCBE0-1000019F	477000BC	94FE3185	05071000	01A54770	0C009603	31890201	01B002B8	.....	.....	-000BEC
CCCCC00-47F00172	5080318C	4180018E	50803174	180058F0	30A405EF	9108318C	471002FC	*..CJ..E.....J..E.....C.....F.....K.....	.....	-000C0C

Figure 4-5. RPG Dump Analysis Listing (Part 9 of 9)



# Appendix A

## Program Exceptions

**Table A-1. Program Exceptions**

Interrupt Code	Interrupt Cause
01	Operation exception: An illegal operation has been attempted or an operation using a noninstalled processor feature has been attempted.
02	Privileged operation exception: A privileged operation has been attempted by a program operating in the problem mode (PS, bit 14 of current PSW, set to 1).
03	Execution exception: The subject instruction of an execute instruction is an execute instruction.
04	Protection exception: A storage protection violation occurs on a program-generated address when the storage protect feature is installed.
05	Address exception: A main storage location outside the range of the installed main storage is referenced by a program-specified address. For the load-control-storage (LCS) instruction only, the referenced control storage location is nonexistent.
06	<p data-bbox="423 1234 643 1262"><u>Specification exception:</u></p> <ul data-bbox="423 1297 1247 1713" style="list-style-type: none"> <li data-bbox="423 1297 1105 1325">• The unit of information referenced is not on an appropriate boundary.</li> <li data-bbox="423 1360 1214 1388">• An invalid modifier field is specified in the service timer register (STR) instruction.</li> <li data-bbox="423 1423 1247 1486">• The <math>r_1</math> field of an instruction that uses an even/odd pair of registers (64-bit operand) does not specify an even register.</li> <li data-bbox="423 1522 1024 1549">• A floating-point register other than 0, 2, 4, or 6 is specified.</li> <li data-bbox="423 1585 1122 1612">• A multiplier or divisor in decimal arithmetic exceeds 15 digits and sign.</li> <li data-bbox="423 1648 1227 1713">• The first operand field is shorter than, or equal in length to, the second operand in decimal, multiply, and divide instructions.</li> </ul>

continued

## Program Exceptions

**Table A-1. Program Exceptions (cont.)**

Interrupt Code	Interrupt Cause
06 (cont.)	<ul style="list-style-type: none"> <li>• The four low-order address bits specified by the contents of <math>r_2</math> comprise a set storage key (SSK) or insert storage key (ISK) instruction and are not equal to 0.</li> <li>• The function specified by the <math>I_2</math> field of a diagnose instruction was not loaded in the transient area of control storage.</li> <li>• A SOFTSCOPE instruction (SSFS or SSRS) was issued without the supporting microcode loaded in the control storage transient area.</li> </ul>
07	<p>Data exception:</p> <ul style="list-style-type: none"> <li>• An invalid sign or digit code is detected in decimal operands.</li> <li>• Fields in decimal arithmetic overlap incorrectly.</li> <li>• The first operand of the multiply decimal instruction does not have sufficient number of high-order 0 digits.</li> </ul>
08	<p>Fixed-point overflow exception: A fixed-point add or subtract operation exceeds the capacity of the first operand field. This interrupt is masked by b, bit 36 of the current PSW.</p>
09	<p>Fixed-point divide exception: The quotient of a fixed-point divide operation exceeds the capacity of the first operand (including division by 0), or the result of a convert-to-binary instruction exceeds 31 bits.</p>
0A	<p>Decimal overflow exception: The result of an add decimal, subtract decimal, or zero-and-add instruction exceeds the capacity of the first operand location. This interrupt is masked by d, bit 37 of the current PSW.</p>
0B	<p>Decimal divide exception: The quotient of a divide decimal (DP) instruction exceeds the capacity of the quotient part of the first operand field.</p>
0C	<p>Exponent overflow exception: The final characteristic resulting from a floating-point arithmetic operation exceeds 127.</p>
0D	<p>Exponent underflow exception: The final characteristic resulting from a floating-point arithmetic operation is less than 0. This interrupt is masked by e, bit 38 of the current PSW.</p>
0E	<p>Significance exception: The final fraction resulting from a floating-point addition or subtraction is equal to 0. This interrupt is masked by s, bit 39 of the current PSW.</p>
0F	<p>Floating-point divide exception: The divisor fraction in a floating-point divide operation is equal to 0.</p>



# Appendix B

## SYSDUMP File Allocation

This appendix contains information about \$YSDUMP file allocation.

Table B-1 shows the number of cylinders required, depending on your system's main storage capacity and the type of disk device you are using.

For example, a system with 8MB main storage requires 53 cylinders for the SYSDUMP file on an 8433 disk drive.

**Table B-1. \$YSDUMP File Size in Cylinders**

Storage Capacity (MB)	Disk Type							
	8416	8417	8418	8419	8430	8433	8470	8494
	Cylinders Required for Use							
1	15	5	15	12	8	7	2	3
1.5	22	8	22	18	11	10	2	4
2	30	10	30	24	15	14	3	5
2.5	37	13	37	30	18	17	4	6
3	44	15	44	36	22	20	4	7
3.5	52	18	52	41	26	23	5	8
4	59	20	59	47	29	27	6	9
5	74	25	74	59	36	33	7	11
6	88	30	88	71	44	40	8	13
7	103	35	103	82	51	46	10	15
8	118	40	118	94	58	53	11	17

---

Use the following procedure to expand the size of the DUMP file after an increase in storage:

1. Use the following JCL to scratch the \$Y\$DUMP file on SYSRES:

```
// JOB SCR DUMP
// DVC 20 // LFD PRNTR
// DVC RES // LBL $Y$DUMP // LFD SYSDUMP
// SCR SYSDUMP
/&
// FIN
```

2. Use Table B-1 to determine the number of cylinders required for your system.
3. Use the following JCL to allocate a MIRAM file (labeled \$Y\$DUMP) and to execute the SG\$OPN load module:

```
// JOB ALLOCATE
// DVC 20 // LFD PRNTR
// DVC RES
// EXT MI,C,1,CYL,xx          (where xx is the number of cylinders)
// LBL $Y$DUMP // LFD SYSDUMP
// EXEC SG$OPN
// PARAM SYSDUMP,100
/&
// FIN
```

You now have sufficient contiguous free space to store a dump of your system's main storage.

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