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**MODELS 92200/93200 SERIES
TIME-MULTIPLEXED
COMMUNICATION CHANNELS
(TMCC)
Technical Manual**

SDS 900685C

September 1965

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SECTION I GENERAL DESCRIPTION

1.1 GENERAL

1.2 This publication provides information relating to the Time-Multiplexed Communication Channel option manufactured by Scientific Data Systems, 1649 Seventeenth Street, Santa Monica, California.

1.3 In this publication, the Time-Multiplexed Communication Channel option is referred to as "TMCC". The models covered, with their description and the figure references, are listed in table 1-1.

Table 1-1. TMCC Models

Model	Description	Channel Used	Figure
92200	6-bit characters with interlace	W(A), C	1-1
92210	6-bit characters with interlace	Y(B), D	1-1
92201	12-bit characters with interlace	C	1-1
92211	12-bit characters with interlace	D	1-1
92202	24-bit characters with interlace	C	1-1
92212	24-bit characters with interlace	D	1-1
93200	6-bit characters without interlace, single channel	W(A), C	1-2
91210	Interlace option for Model 93200		
93201	12-bit character extension option for Model 93200		
93202	24-bit character extension option for Model 93200		
93221	6-bit characters without interlace, two channels	W + Y (A + B)	1-2
91210	Interlace option for either channel of Model 93221		
93201	12-bit character extension option for Model 93221		
93202	24-bit character extension option for Model 93221		

1.4 The information in this publication relates to the TMCC as utilized with the 925/930/9300 computers. Other publications containing information relating to the TMCC and input/output operation are listed in table 1-2.

Table 1-2. Applicable Publications

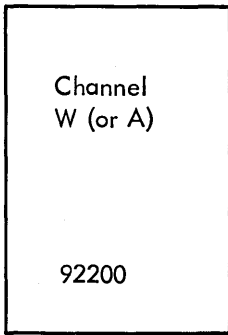
Title of Publication	Publication No.
SDS 925/930/9300 TMCC Input/Output Unit Logic Layouts, Current and History	900557
SDS 925 Computer Reference Manual	900099
Model 925 Computer, Technical Manual	900633
SDS 930 Computer Reference Manual	900064
Model 930 Computer, Technical Manual	900066
SDS 9300 Computer Reference Manual	900050
Model 9300 Computer, Technical Manual	900570

1.5 PURPOSE OF OPTION

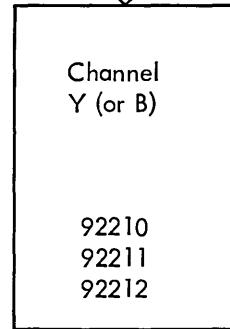
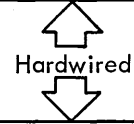
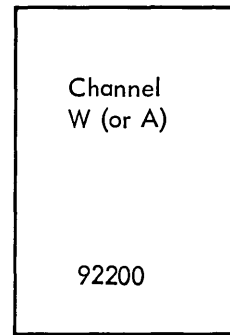
1.6 The TMCC is a time-multiplexed input/output channel utilized for communication between peripheral devices and the 925/930/9300 computers. Its operation is designated "time-multiplexed" because it gains access to the computer memory through the same path utilized by the computer and must, therefore, momentarily interrupt computation to store or obtain a word of information. Up to four TMCCs may be connected to one computer and all may be active simultaneously but since their operation is time-multiplexed, only one channel at a time communicates with the computer memory.

1.7 DESCRIPTION AND LEADING PARTICULARS

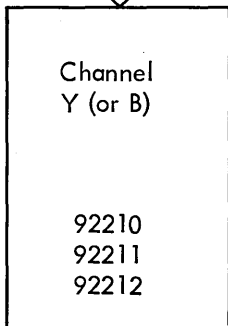
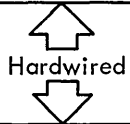
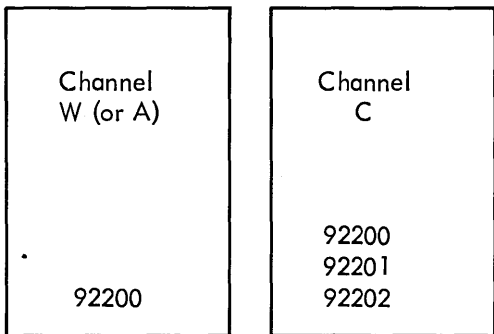
1.8 A computer may have from one to four TMCCs connected to it. These are designated by letter symbols in the order of their installation. When only one



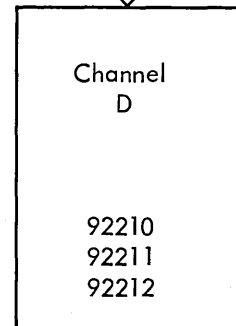
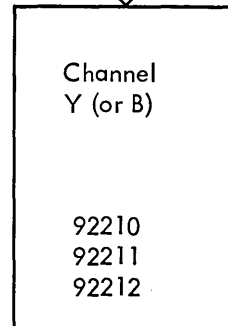
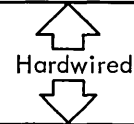
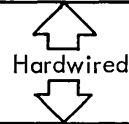
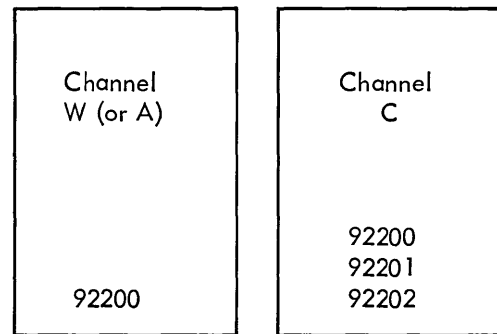
One Channel Configuration



Two Channel Configuration

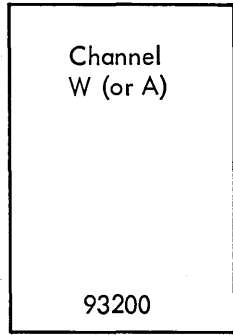


Three Channel Configuration

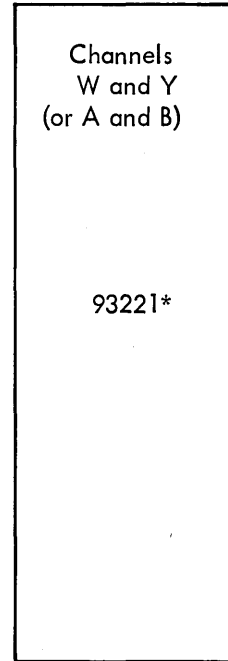


Four Channel Configuration

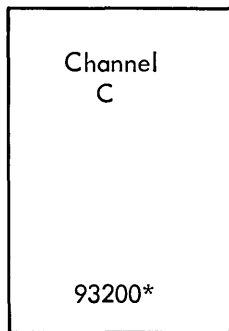
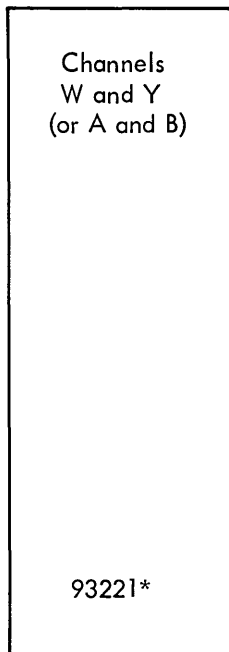
Figure 1-1. TMCC Configuration, Models 922XX



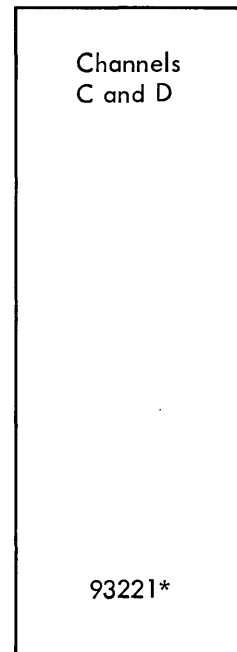
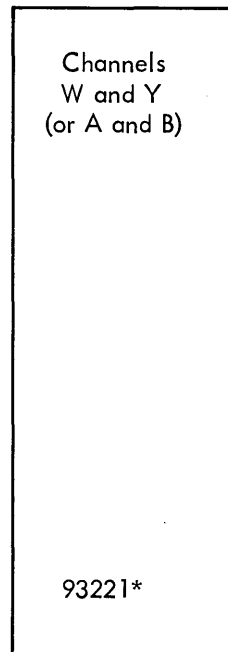
One Channel Configuration



Two Channel Configuration



Three Channel Configuration



Four Channel Configuration

*Note: The Interlace feature, Model 91210, is required on channels B, C and D, and is optional on W, A and Y. All channels may have 93201 or 93202 options.

Figure 1-2. TMCC Configuration, Models 932XX

channel is used it is specified as the W channel on the 925 and 930 computers or as the A channel on the 9300 computer. As additional channels are added they are designated, in order, Y (or B on the 9300 computer), C, and D. Because single-word input/output instructions (Memory Into W, Memory Into Y, Memory Into A, W Into Memory, Y Into Memory, and A Into Memory) are not available for all channels, the interlace feature is mandatory on channels B, C, and D.

1.9 Primary differences between the models occur in the hardware layouts, connections between channels, and the options available.

1-10 PHYSICAL DESCRIPTION

1.11 The TMCC consists of plug-in modules contained in chassis consisting of four or six rows of modules.

Each row contains 32 connectors thereby allowing the insertion of up to 32 modules in each row. Physical location of each module is given in Section 4 of this manual.

1.12 TMCC Models 922XX

1.13 In TMCC Models 922XX, the character length is fixed for a particular model and the interlace feature is included. One channel must be six bits to allow entering a program into the computer.

1.14 One Channel Configuration

1.15 A one channel configuration TMCC (Model 92200) consists of a single chassis containing four rows of modules (C through F). The quantity and type of modules are listed in table 1-3.

Table 1-3. Models 922XX Module Complement

Item	Description	Model	Quantity					
			92200 6-Bit "W" "A" or "C"	92201 12-Bit	92202 24-Bit	92210 6-Bit "Y" "B" or "D"	92211 12-Bit "Y" "B" or "D"	92212 24-Bit "Y" "B" or "D"
1	Triple Flip-Flop	FB52	19	21	25	18	20	24
2	NAND Flip-Flop	FB54	14	14	14	13	13	13
3	Cable Driver No. 2	AB55	2	2	2	1		
4	NAND No. 2	IB56	5	5	5	5	5	5
5	Band NAND	IB57	4	4	4	4	4	4
6	NAND Module	IB52	2	2	2			
7	Shift Register	DB50	3	3	3	3	3	3
8	Receiver Inverter Buffer	AB53	2	2	2			
9	Termination Module	ZB52	2	2	2			
10	Interface +8 to +4	NB50	1	2	2	1	1	1
11	Cable Driver	AK53	9	10	11	3	1	2
12	Cable Driver	NB52	3	3	3			
13	Termination Module +4	ZB50	7	7	7	7	7	7
14	Termination Module	ZB55	1			1		
15	Receiver, Inverter	AB52	2	2	2			
16	NAND No. 4	IB59	5	5	5	4	4	4
17	Schmitt Trigger	AK54	1	1	1			
18	Termination Module	ZB56		1			2	
19	Termination Module	ZB57			1			2

1.16 Two Channel Configuration

1.17 A two channel configuration TMCC consists of two chassis, one containing four rows of modules (C through F) and the other containing two rows of modules (A and B). These two chassis are physically bolted together and hardwired to allow mounting as a single unit. The two channel configuration consists of a Model 92200 for channel W (or A) and either a six bit option (92210), a twelve bit option (92211) or a twenty-four bit option (92212) for channel Y (B). The options selected determine the quantity and type of modules. Table 1-3 lists the modules required for each model.

1.18 Three Channel Configuration

1.19 The three channel configuration consists of three chassis, two containing four rows of modules and the third containing two rows of modules. Channel W (or A), consists of a Model 92200; channel Y (or B) consists of a six bit option (92210), a twelve bit option (92211), or a twenty-four bit option (92212); channel C consists of a six bit option (92200), a twelve bit option (92201), or a twenty-four bit option (92202). The quantity and types of modules for each option are listed in table 1-3.

1.20 Four Channel Configuration

1.21 The four channel configuration consists of four chassis, two containing four and two containing two rows of modules. The four channel configuration consists of a Model 92200 for channel W (A); a six bit option (92210), a twelve bit option (92211) or a twenty-four bit option (92212) for channel Y (B); a six bit option (92200), a twelve bit option (92201), or twenty-four bit option (92202) for channel C; and a six bit option (92210), twelve bit option (92211), or a twenty-four bit option (92212) for channel D. The quantity and types of modules are listed in table 1-3.

1.22 TMCC Models 932XX

1.23 In TMCC Models 932XX, the interlace registers are optional equipment. Interlace options may be installed in any of the four TMCCs that a 925/930/9300 computer may have. However, channels C and D (and B on the 9300 computer) must have interlace installed since there are no computer instructions allowing use of these channels without it.

1.24 One Channel Configuration

1.25 A one channel configuration TMCC (Model 93200) consists of a single chassis containing four rows of modules. The quantity and types of modules are listed in table 1-4.

1.26 Two Channel Configuration

1.27 The two channel configuration consists of a single chassis containing six rows of modules. For channels W and Y, a Model 93221 TMCC may be used. The Model 93221 may be extended to twelve bit or twenty-four bit characters by addition of modules as listed in table 1-4. The Model 93221 must have the Model 91210 Interlace installed in either or both channels. The modules comprising the Model 91210 Interlace are indicated in table 1-4.

1.28 Three Channel Configuration

1.29 A three channel configuration TMCC consists of a single chassis containing six rows of modules and another chassis of four rows of modules. Channels W and Y (or A and B) are as given in paragraph 1.27. Channel C consists of a Model 93200 TMCC (with character extension options) and must include a Model 91210 Interlace.

1.30 Four Channel Configuration

1.31 A four channel configuration TMCC consists of two chassis, each containing six rows of modules. Two Model 93221 TMCCs may be used with character extension options if desired. Model 91210 Interlace options must be installed on at least channels C, and D (and B on the 9300).

1.32 SEMICONDUCTOR COMPLEMENT

1.33 The semiconductor complement may be derived from the Material Lists and Module Data Sheets contained in Section 6. The module complement for each model of TMCC is contained in tables 1-3 and 1-4.

1.34 INTERLACE FEATURE

1.35 The purpose of the interlace feature is to provide the TMCC with a means of transferring blocks of words without requiring a separate instruction for each word. To do this, two counters are added to the TMCC. One counter is loaded with a count of the number of words in the block, and the other counter is loaded with the address of the memory position of the first word in the block. Then, as each input or output word is transferred to or from memory, the Word Counter and the Address Counter are incremented. The Word Counter holds the one's complement of the count; thus the count is decreased as the counter counts up. In addition to the two counters, the interlace logic also includes Channel Command Interrupt Enables, and all the necessary control logic.

Table 1-4. Models 923XX Module Complement

Item	Description	Model	Quantity				
			93200 6-Bit "W"	93221 6-Bit "W" & "Y"	93201 12-Bit "W" or "Y"	93202 24-Bit "W" or "Y"	91210 Interlace "W" or "Y"
1	Triple Flip-Flop	FB52	6	11	2	6	13
2	NAND Flip-Flop	FB54	13	25	-	-	1
3	Cable Driver No. 2*	AB55	3	4	-	-	-
4	NAND No. 2	IB56	4	8	-	-	1
5	Band NAND	IB57	3	7	-	-	1
6	NAND Module	IB52	2	2	-	-	-
7	Shift Register	DB50-2	3	6	-	-	-
8	Receiver Inverter Buffer	AB53	2	2	-	-	-
9	Termination Module	ZB52	2	2	-	-	-
10	Interface +8 to +4	NB50	1	2	1	1	-
11	Cable Driver	AK53	9	13	1	2	-
12	Driver Cable Interface	NB52	1	1	-	-	2
13	Termination Module +4	ZB50	6	10	0	0	1
14	Priority Interrupt**	SK61	2	2	-	-	-
15	Receiver Inverter	AB52	2	2	-	-	-
16	NAND No. 4	IB59	3	5	-	-	3
17	Schmitt Trigger	AK54	1	1	-	-	-

*SK61s are located in computer basic interrupt: 925/930 Reference Drawing 107352
9300 Reference Drawing 107626

**One AB55 is located in 5E (930 only) if I/O is used as a "C" or "C-D" channel.

SECTION II PROGRAMMING

2.1 PURPOSE

2.2 The 925/930/9300 computers include as standard equipment one Time-Multiplexed Communication Channel (TMCC), without interlacing capability, as well as provision for three additional channels. The interlace unit is available as an option. The W and Y channels are available with or without interlace; the C and D channels are available only with interlace. The W channel on the 925 and 930 computers is equivalent to the A channel on the 9300 computer and the Y channel on the 925 and 930 computers is equivalent to the B channel on the 9300 computer. (The B channel on the 9300 must also have the interlace feature). These channels are capable of automatically controlling the flow of data to and from memory at rates up to one word every 3.5 microseconds. These channels run independently of the central processor and only communicate with it to transfer data to or from memory.

2.3 GENERAL OPERATION

2.4 Utilizing channels W and Y (or A), characters, and words can be transmitted between memory and peripheral devices under the direct control of single instructions. Each of these channels has associated with it two instructions to facilitate direct control operations. For channel W, W INTO MEMORY (WIM) (channel A, A INTO MEMORY (AIM)), causes a word from a peripheral transmission to be taken from the channel W (A) buffer register and placed directly in the specified memory location without disturbing any internal registers. MEMORY INTO W (MIW), (MEMORY INTO A (MIA)), causes a word to be taken from a specified memory location and placed in the channel W (A) buffer register to be read out to the currently operating peripheral device connected to the channel. WIM (AIM) and MIW (MIA) are preceded by instructions from the EOM group that set up the input/output operation. YIM and MIY instructions function in an analogous manner for channel Y. The general test instruction, SKIP IF SIGNAL NOT SET (SKS) provides the facility for testing error indications and/or for testing various peripheral device indicators.

2.5 Additionally, using any channel including channels W and Y (A and B) with interlace, data can be transmitted to and from core storage under channel control. Operation of a channel is initiated by the execution of a sequence of instructions in the central processor. Once started, the channel operates independently of the central processor, automatically transferring each word at the correct time.

2.6 Three instructions control the process of transmitting and receiving data between channel peripheral equipment and the central processor. These instructions are:

EOM	ENERGIZE OUTPUT M
POT	PARALLEL OUTPUT
SKS	SKIP IF SIGNAL NOT SET

2.7 EOM instructions activate one of channels W (A), Y (B), C, or D, to select the peripheral device to be used, and to set up the initial conditions of the data transmission, including the peripheral operation to be performed. An EOM instruction also specifies terminal conditions for an operation.

2.8 PARALLEL OUTPUT (POT) instruction sends out to the channel the number of words in the transmission and the address at which the output begins.

2.9 SKIP IF SIGNAL NOT SET (SKS) instruction can test the Error indicators, End-of-Transmission indicators, and other input/output control indicators, such as printer end-of-form or card hopper empty.

2.10 The general order of use of these instructions for interlaced operation is:

<u>Instruction</u>	<u>Function</u>
EOM	to address the channel, connect the peripheral device, specify various input/output conditions, and alert the optional channel interlace (see Communication Channel Input/Output, paragraph 2.30).
EOM	to specify the terminal conditions and interrupts desired during the transmission
POT	to transmit to the channel a word containing the transmission starting address and block length

Bits 0 through 9 of this latter word contain the ten lower order bits of the word count; bits 10 through 23 contain the 14 bits of the starting address. The second EOM contains the high-order bits of the word count and starting address when needed.

2.11 DIRECT PARALLEL INPUT/OUTPUT

2.12 The direct parallel input/output (POT/PIN) facility allows any word in core memory to be presented, in

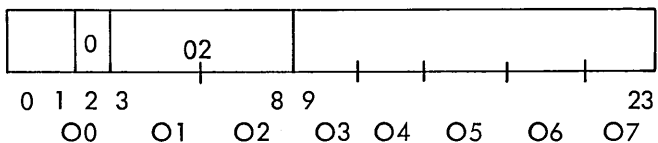
parallel, at any special system connector or applicable standard peripheral connector; or, conversely, allows signals sent to a connector to be stored in any core memory location. EOM and SKS instructions control parallel input/output operations in the same way as in channel operations. POT/PIN instructions also generate or check for correct parity with each word transmitted.

2.13 SINGLE-BIT INPUT/OUTPUT

2.14 EOM and SKS instructions also perform single-bit input/output and testing for special or standard devices. The execution of an EOM transmits a single signal of approximately 1.4 microseconds duration to an external connector and also provides the connector with a 15-bit address for the destination of this signal. SKS tests whether a similar signal is present on an external connector and skips accordingly.

2.15 PRIMARY INPUT/OUTPUT INSTRUCTIONS

2.16 ENERGIZE OUTPUT M (EOM)



2.17 The major instruction for preparing channel W (or Y, C, D) and an attached peripheral device to perform a data transmission or other peripheral activity is the multi-purpose instruction, ENERGIZE OUTPUT M (EOM). This instruction operates in four distinct modes with many functional configurations. These modes are Buffer Control, Input/Output Control, Internal Control, and System Control. In the third and fourth modes, EOM controls and initiates non-communication channel operations such as special systems transmissions. Each of the frequently used EOM instruction configurations has a mnemonic tag used with standard SDS assemblers. The different modes of operation are program-selectable by the setting of two bits (10, 11 of octal position 3) within the EOM instruction format:

Octal Value	Bit Position 10	Bit Position 11	Area
0	0	0	Buffer Control
1	0	1	Input/Output Control
2	1	0	Internal Control
3	1	1	System Control

2.18 A Buffer Control mode EOM operates essentially as a set-up or preparation facility for data transmissions

or other peripheral activities using the channel. The channel to be used, the peripheral unit on that channel, the operation to be performed, and the type of character format to be used are all detailed within this EOM. It also details the use of BCD or binary data transmission, the allowance or not of a leader (as in paper tape), and the direction of operation (as in forward direction for magnetic tape). Execution of such an EOM "connects" the specified peripheral unit to the channel. An EOM in this mode can also alert the interlace, which is the optional, automatic buffer control for input/output.

2.19 An EOM in the Input/Output mode directs peripheral devices to perform non-transmitting operations such as rewind magnetic tape and upspace the printer. This EOM selects certain channel operations such as interrupt response and input/output terminal function desired. It alerts peripheral devices that a PARALLEL INPUT (PIN) or PARALLEL OUTPUT (POT) instruction follows. It also can give an extension of the word count to 15-bits for the number of words to be transmitted and an extension of the address specification to 15-bits. Without disturbing the associated channel, this EOM can also set up the interlace unit. It is with the input/output mode EOM that the user selects his I/O operation as compatible or extended I/O modes.

2.20 This coding sequence initiates such an interlaced channel operation (compatible mode):

Instruction	Function
EOM (Input/Output Control Mode)	Alert the interlace
POT	transmit starting address and block length to interlace
EOM (Buffer Control Mode)	address channel, connect peripheral device, specify various input/output conditions, start transmission

2.21 Initiating an interlaced input/output operation via this sequence of instructions facilitates checkout by allowing the programmer to single-step through this portion of the program. The first two instructions, EOM (loc) and POT, set up the interlace with data address and block length. Therefore, single-stepping through the sequence allows the interlaced channel to complete the input/output operation. When a single EOM (Buffer Control mode) sets up the channel and interlace with a POT instruction following, the programmer cannot step through the sequence since the input/output operation proceeds before the next stepped instruction (POT) places the address and block length in the interlace.

2.22 An EOM in the Internal Control mode enables and disables the interrupt system. EOM in this mode also

End-of-Record and Zero Word Count, come from the interlace and are under its control. The interlace controls input/output termination functions during interlaced operation.

2.38 Two EOM instructions and a POT instruction alert and set up a channel interlace. The first EOM alerts the interlace, that is activates the interlace and instructs it to expect a word count and starting address to be sent to it by the POT instruction. The second EOM is an Input/Output mode EOM that specifies the interrupt and the terminal function to be used. This EOM also can specify a 15th address bit and five more high-order word count bits expanding the word count from 10-bits to 15. This sequence is written: EOM (Alert), EOM (I/O), and POT. When the channel buffer is being set up at the same time, the buffer control EOM can alert the interlace. When the buffer is already set up, during a continuing I/O operation, the programmer may use the I/O EOM, ALERT CHANNEL (00250000), to alert the interlace.

2.39 When the programmer does not desire to program the Extended Mode with the input/output terminal functions, interrupts, and additional count or address, only the EOM (Alert) and the POT are necessary to set up the channel interlace (Compatible mode).

2.40 In the Extended Mode, the four channels are programmed in the same way.

2.41 The Time-Multiplexed Channels use the memory logic of the central processor to facilitate input and output of data words. The transfer of each word between a time-multiplexed channel buffer and memory requires two memory cycles. During this time, computation stops in the central processor. Priority for the use of the word input/output logic is in the order: Channel D, C, Y(B), W(A). Any Time-Multiplexed Channel operating with interlace has priority over the central processor for memory access.

2.42 COMMUNICATION CHANNEL DESCRIPTION

2.43 Up to 30 peripheral devices may be connected to one channel. Each of these devices has a unique, two-digit, octal address by which it is selected for an input/output operation. To select the peripheral device, the program loads the proper unit address into the 6-bit Unit Address Register (UAR) in the channel buffer. This address selects both the device and, if appropriate, the function to be performed. Placing a non-zero unit address in the Unit Address Register "connects" the peripheral unit addressed to the channel and it becomes "active". When the UAR contains a zero address, or any time that a terminal or initial condition clears the contents of UAR, the channel is "inactive." The zero

in UAR also means that it is not connected to a peripheral unit.

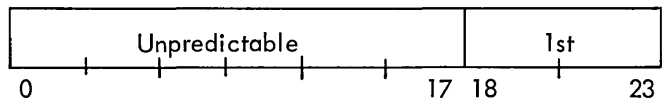
2.44 When the channel and the peripheral unit to be used have been connected, the channel must have information pertaining to the location in memory of the data to be transmitted or received and pertaining to the number of data words in the transfer.

2.45 TIME-MULTIPLEXED CHANNEL REGISTERS

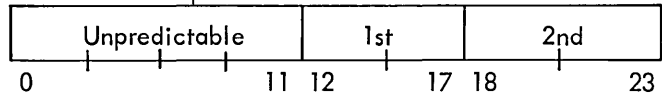
2.46 In the Time-Multiplexed Channels W (A) through D, there are two registers important to the programmer, the Word Assembly Register (WAR) and the Single-Character Register (SCR). The WAR, a 24-bit, word-sized buffer, contains the word of data actively being received or transmitted during an input or output operation. During input, 6-bit characters (plus parity) enter the Single-Character Register where the channel buffer assembles them, one at a time, into the WAR. Then the completed word is placed in memory. Depending on the number of characters per word specified, the word assembled and placed in memory during input has the form:

Word in Memory

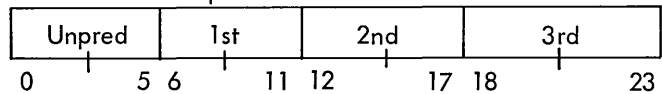
One character per word mode



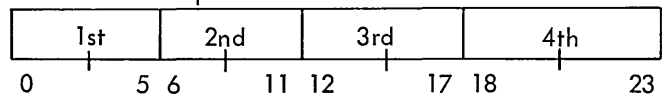
Two characters per word mode



Three characters per word mode



Four characters per word mode



2.47 The unfilled character positions contain unpredictable data. When assembled during a one-word operation, a WIM (AIM) instruction places the word into memory. Under interlace control, the interlaced channel automatically places the word in memory when assembled.

2.48 When the end of an information record is detected by a buffer, the buffer automatically disengages from

the device and is then "ready" for another operation. The buffer logic is reset, except that the state of the error indicator is maintained and the last word of the input is still in the word register. If the number of characters in the input record was not a multiple of the number of characters assembled into each computer word, then zeros are automatically forced into the least significant positions of the last word. This last word can then be stored in memory by a BUFFER INTO M WHEN READY WIM (AIM) or YIM instruction after the buffer has disengaged. If the number of characters in the input record was a multiple of the number of characters assembled into each computer word, then the word remaining in the W buffer is either the last group of characters from the input device, if they were not previously transferred to memory by a BUFFER INTO M WHEN READY WIM (AIM) or YIM, or zeros if the last group of characters had been transferred to memory. In either case, it is safe to issue one such instruction after the buffer has disengaged without "hanging up" the computer.

2.49 During output, words come from memory into the WAR where the channel buffer disassembles them into the SCR one 6-bit character at a time. Depending on the characters per word mode specified, the 6-bit characters within the word are output as follows:

<u>Function</u>	<u>Mode</u>
Output one character from bits 0 through 5	One character per word
Output two characters from bits 0 through 5, 6 through 11	Two characters per word
Output three characters from bits 0 through 5, 6 through 11, 12 through 17	Three characters per word
Output four characters from bits 0 through 5, 6 through 11, 12 through 17, 18 through 23	Four characters per word

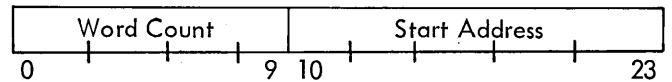
2.50 As required, the characters are transferred into the One Character Register and output with generated parity. After each character transfer, the word in the WAR is shifted left six bits to be ready for the next transfer. Only those characters needed from each word are used; when required, a new word is brought to the WAR for the next character. For special applications a Time-Multiplexed Channel may be equipped with a

12- or 24-bit One Character Register. The external device which has a character size greater than 6-bits specifies to the channel what its size is, 12- or 24-bits. Standard 6-bit devices are unaffected by the installation of a wider SCR.

2.51 Interlace Registers

2.52 A channel interlace contains two working registers, the Word Count Register (WCR) and the Memory Address Register (MAR). In the set-up sequence -- EOM, EOM, POT -- for an interlaced input/output operation, the POT instruction transmits to the interlace a data word made up of the word count (that is, length) and the starting address of the data block. The 15-bit Word Count Register (WCR) contains the data word count during a data transfer. The number of data words is decremented by one and the new count replaces the old one in the WCR for each word transmitted.

2.53 The count is assembled into the WCR from two places: the least significant 10-bits is from the "POTed" word and the most significant 5-bits is from the "HI COUNT" field of the second EOM. The form of the "POTed" word is:

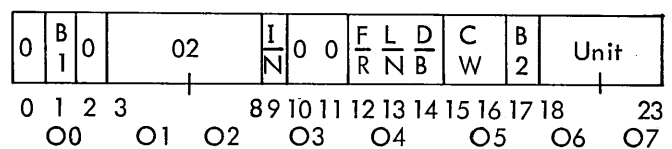


2.54 When the word count is equal to zero, the transmission is complete. During output, this causes a termination; during input, the interlace allows any further data to fill the channel buffer and generates the End-of-Word interrupt, if enabled.

2.55 The Memory Address Register (MAR) contains the starting destination or source address in memory of the transmitted data. The memory locations to or from which data words are to be transmitted enter the MAR at the same time the word count does. During transmission of data, the interlace increments the contents of the MAR after each word as it decrements the contents of the WCR. These two registers provide the interlace control of block transmissions. The high-order 15th address bit comes from the second EOM, also.

2.56 COMMUNICATION CHANNEL PROGRAMMING

2.57 The ENERGIZE OUTPUT M (EOM) used in the Buffer Control mode addresses and connects the specified Channel W (A), Y (B), C, or D, and selects the desired unit address. The detailed instruction format is:

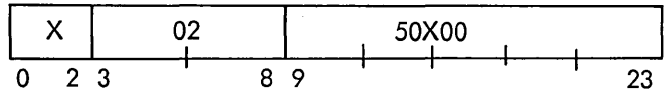


<u>Bit Designation</u>	<u>Octal Position</u>	<u>Octal Value</u>	<u>Function</u>
B1	O0	2	Bit positions 1 and 17 specify the channel to be activated.
B2	O5	1	Channel W (A) is numbered 00, channel Y (B) is 01, channel C is 10, and channel D is 11.
O2	O1-2	02	Bit positions 3 through 8 contain 02, the instruction code for EOM.
I/N	O3	4	A 1-bit in position 9 alerts the buffer interlace.
00	O3	0	Bit positions 10 and 11 contain the EOM mode indicator for the Buffer Control mode.
F/R	O4	4	Bit position 12 specifies the direction in which the peripheral device will operate. A "0" specifies the forward direction. A "1" specifies the reverse direction.
L/N	O4	2	Bit position 13 specifies whether the device should be started with a leader as in paper tape. A "0" specifies a start with leader. A "1" specifies a start without leader.
D/B	O4	1	Bit position 14 specifies the mode of character format. A "0" specifies BCD format. A "1" specifies Binary format.
C/W	O5		Bit positions 15 and 16 specify the number of characters to be assembled into, or disassembled from, each transmitted word. One character per word is specified by 00, two by 01, three by 10 and four by 11.
UNIT	O6-7		Bit positions 18 through 23 specify the unit and the function to be performed with that unit.

2.58 STANDARD EOM CHANNEL INSTRUCTIONS

2.59 Several EOM function configurations have standard uses. These have standard, assembler-type mnemonics and are separate instructions.

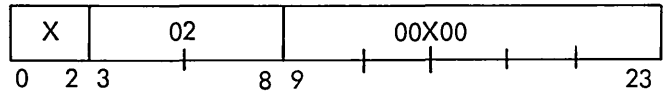
2.60 ALERT CHANNEL (ALC)



2.61 ALC alerts the channel interlace. This instruction does not disturb the channel buffer in any way. ALC has no effect on W or Y Buffers without interlace. The channel Alerts are:

<u>Mnemonic</u>	<u>Alert Channel</u>	<u>Instruction</u>
ALC 0	W (A)	0 02 50000
ALC 1	Y (B)	0 02 50100
ALC 2	C	2 02 50000
ALC 3	D	2 02 50100

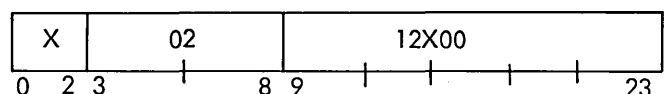
2.62 DISCONNECT CHANNEL (DSC)



2.63 DSC disconnects the channel. It unconditionally sets the Unit Address Register to 00 regardless of whether the channel is currently addressing a device. This instruction disconnects any device which may be connected to the channel. It also unconditionally makes the channel Ready (Inactive) and clears the Error indicator.

<u>Mnemonic</u>	<u>Disconnect Channel</u>	<u>Instruction</u>
DSC 0	W (A)	0 02 00000
DSC 1	Y (B)	0 02 00100
DSC 2	C	2 02 00000
DSC 3	D	2 02 00100

2.64 ALERT TO STORE ADDRESS FROM CHANNEL (ASC)



2.65 ASC alerts an interlaced channel so the PIN instruction that follows can store the contents of the Memory Address Register. This instruction does not affect the operation of the channel.

2.66 ASC is always used in conjunction with PIN to determine the current status of a peripheral operation being performed by the selected channel. The two instructions are written together:

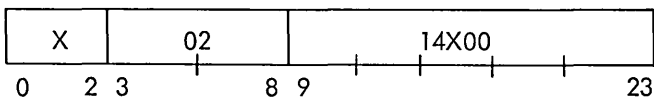
ASC n
PIN m, x

2.67 When the program executes these two instructions, the contents of the effective memory location designated by the PIN instruction are:

Bit Positions	Contents
0 through 8	Zero
9 through 23	Contents of channel's Memory Address Register

Mnemonic	Channel	Instruction
ASC 0	W (A)	0 02 12000
ASC 1	Y (B)	0 02 12100
ASC 2	C	2 02 12000
ASC 3	D	2 02 12100

2.68 TERMINATE OUTPUT OF CHANNEL (TOP)



2.69 When the last word of a block enters the channel, TOP terminates channel output. After the execution of this instruction, the following occurs. When the channel buffer delivers the last character to the peripheral device, the buffer disconnects.

2.70 TOP always terminates a non-interlaced channel output operation. It may be used with all communication channels if the particular function selected is terminal function 11 but no further data output is required.

Mnemonic	Terminate Output on Channel	Instruction
TOP 0	W (A)	0 02 14000
TOP 1	Y (B)	0 02 14100
TOP 2	C	2 02 14000
TOP 3	D	2 02 14100

2.71 COMPATIBLE/EXTENDED INPUT/OUTPUT MODES

2.72 The termination of an I/O operation and the interrupts that may be associated with that termination fall into two classes: Compatible and Extended. The choice of one of these two "modes" of input/output operation determines how the system behaves when the termination of an I/O operation occurs.

2.73 Interrupts occurring at the same level (e.g., location 30, 31, etc.) can have difference names (e.g., Count Equal Zero and End-of-Word). These names reflect the different I/O mode in operation when the interrupt occurs. The differences include the timing of interrupt occurrence relative to the I/O operation and type of interrupt requested.

2.74 The Compatible mode of operation for channels W(A), Y(B), C, D is directly compatible with the SDS 900 Computer series modes of I/O operation. The types of interrupts that can be requested are the End-of-Word and End-of-Transmission interrupts.

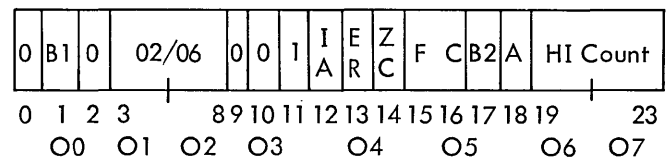
2.75 The Extended mode for all channels expands the I/O capabilities to include the "terminal functions" discussed below. The types of interrupts that can be requested are the Count Equal Zero and End-of-Record interrupts.

2.76 The I/O mode is selected in the Input/Output EOM via bit 12, the Interrupt Arm bit. A 0-bit makes the system operate in the Compatible mode; a 1-bit sets the system in the Extended mode.

2.77 In particular, the Interrupt Arm (IA) bit determines whether any of the Extended functions operate; that is, a "0" in IA means that the other Extended mode controls, bits 13, 14, 15 and 16, have no effect.

2.78 INPUT/OUTPUT CLASS EOM

2.79 The Input/Output EOM selects the I/O operation mode. When the Extended mode is selected, this EOM also selects (arms) which interrupts are to be operational and selects the desired terminal function. This EOM applies to channels W(A), Y(B), C, and D.



<u>Bit Designation</u>	<u>Octal Position</u>	<u>Octal Value</u>	<u>Function</u>
0			Bit positions 0 and 2 are not used with this EOM.
B1 B2	O0 O5	2 1	Bit positions 1 and 17 specify the channel.
02/06	O1-O2	02	Bit positions 3 through 8 contain 02/06, the instruction code for EOM.
01	O3	1	Bit positions 10 and 11 contain the EOM indicator for the Input/Output control mode.
IA	O4	4	Bit position 12 selects the mode of I/O operation. A "0" specifies the Compatible mode. The operation of bits 13, 14, 15, and 16 are disallowed. Channels W (A), Y (B), C and D operate in this mode which is completely SDS 900 series compatible. If interrupts are required, the user enables the Interrupt System, thus enabling and arming the End-of-Word and End-of-Transmission interrupts. A "1" specifies the Extended mode. All channels can operate in this mode. This allows the use of bits 13, 14, 15, and 16. If interrupts are required, the user arms the associated ones by placing 1-bit in bits 13 and/or 14. The "terminal function" to be used is selected via bits 15 and 16.
ER		2	Bit position 13 controls the arming of the End-of-Record interrupt. A 1-bit arms the interrupt. A 0-bit disarms the interrupt.

<u>Bit Designation</u>	<u>Octal Position</u>	<u>Octal Value</u>	<u>Function</u>
ZC		1	Bit position 14 controls the arming of the Zero Word Count interrupt. A 1-bit arms the interrupt. A 0-bit disarms the interrupt.
FC	O5		Bit positions 15 and 16 specify the terminal condition function to be performed with the transmission.
A			Bit position 18 is the high-order address bit.
HI Count			Bit positions 19 through 23 contain the most significant four bits of the 15-bit word count. These positions specify a word count greater than 1023.

NOTE

A 1-bit in 13 and/or 14 does the following:

1. Arms that interrupt during this complete I/O operation; disconnecting this channel disarms the interrupt.
2. Once armed by bits 13 and/or 14, the interrupt can be enabled or disabled by the Enable/Disable feature of the Interrupt System. If a channel generates an extended mode I/O interrupt while the system is disabled, the designated interrupt level goes to the Waiting state. When the program again enables the interrupt system, the interrupt goes to the Active state when its priority allows.

2.80 TERMINAL FUNCTIONS; EXTENDED MODE

2.81 A 2-bit function code in the Input/Output EOM controls the termination of input/output operation in the extended mode. These functions are described below with the letter C representing the specified word count of the transmission.

2.82 INPUT/OUTPUT OF A RECORD AND DISCONNECT (IORD)

2.83 Input

2.84 Read C words. If C equals zero before the End-of-Record is detected, the rest of the record is ignored. At the End-of-Record, the peripheral device is disconnected and the channel becomes inactive.

2.85 Output

2.86 Write C words. When C equals zero, output is terminated (i.e., the device is signaled that the last characters have been transmitted). When the peripheral device has generated the end of record and, if necessary, checked the validity of the record, it sends an End-of-Record response to the channel buffer. When received by the buffer, the End-of-Record signal generates an End-of-Record interrupt (if armed) and disconnects the channel.

2.87 The line printer generates the End-of-Record response when it completes the printing of a line. If the printer encounters any print errors or faults, it sends a signal to the channel that sets the channel error indicator; this can occur since the printer has not disconnected from the channel. The IORD is useful when the program is to print several lines and the program is not otherwise to use the channel between lines. When the printer completes each line, it causes an End-of-Record interrupt (assumed to be armed), notifying the program that it can immediately transmit the next paper control instruction and the next line image.

2.88 The unbuffered card punch operates similarly. It generates the End-of-Record response after punching each row. If any faults occur during the punching of the entire card, the card punch sends a signal to the channel that sets the channel error indicator; this occurs after punching the last row (row 9).

NOTE

A program should not use IORD with devices that do not have End-of-Record conditions on input (e.g., typewriter) or generate End-of-Record responses upon output termination, (e.g., devices such as the paper tape punch and typewriter). These devices do terminate output but give the program no indication when they receive the last characters.

2.89 INPUT/OUTPUT UNTIL SIGNAL THEN DISCONNECT (IOSD)

2.90 Input

2.91 Read C words. When C equals zero or when the End-of-Record is encountered, the device is disconnected and the channel becomes inactive. If the channel disconnects because of a zero count, an EOR interrupt (if armed) will be generated in addition to the count equal zero interrupt. If both are armed, C=0 will occur first.

2.92 Output

2.93 Write C words. When C equals zero and when the last character has been transmitted, the channel disconnects the device and becomes inactive. If an End-of-Record signal is received before the count reaches zero, the channel will disconnect immediately.

NOTE

The IOSD is designed for use on devices which are normally operated on the basis of the word count only. Typewriters and paper tape devices are of this type, as are the printer and card punch when the user does not wish to stay connected until the operation is complete.

2.94 INPUT/OUTPUT OF A RECORD AND PROCEED (IORP)

2.95 Input

2.96 Read C words. If the channel counts C down to zero before the peripheral device encounters the End-of-Record (EOR), the channel ignores the rest of the record (to the End-of-Record). When the peripheral device sends the End-of-Record signal to the channel, the channel sets its End-of-Record Indicator; this signal sets the End-of-Record interrupt (if armed). The channel does not disconnect. The channel is now in an "Inter-record" condition.

2.97 When the peripheral device is magnetic tape, the tape continues to move when the tape handler encounters the End-of-Record. The End-of-Record occurs when the tape read-heads encounter tape gap; this also causes a Tape Gap signal to "come high". If the program executes a new read tape or scan tape EOM during the inter-gap time (approximately one millisecond while the Tape Gap signal is high), the tape remains in motion and proceeds to read or scan the next record.

If the program executes no such EOM before the Tape Gap signal drops, the channel disconnects and the tape comes to a stop. No additional interrupt occurs. This is the only condition that causes a channel to disconnect automatically.

2.98 All other input devices remain connected until the program takes further action. The paper tape reader remains in motion; the program should issue a "disconnect channel" instruction if the program is not reading any more tape. To proceed after the End-of-Record occurs, the program first executes a Buffer Control mode EOM to re-initialize the Channel Unit Address Register and then reloads the interlace portion of the channel (the program can alert the Interlace via the Buffer Control EOM). Otherwise, the channel immediately terminates any attempt to use its interlace portion since the channel is aware that it is still active and in the End-of-Record condition. When the program continues from an Inter-record condition, the program should use an extended mode terminal function. An IORP should not be used to read devices which do not have EOR signals (e. g., the typewriter and paper tape punch).

2.99 Output

2.100 Write C words. When the channel interlace counts C down to zero, the Interlace notifies the channel buffer that it has received the last word that is to be output; when the buffer outputs this last word, it sends a signal to the connected peripheral device indicating that the device has the last word now. When the peripheral device "receives, outputs and checks the validity of" this last word, it sends an End-of-Record response to the channel buffer. When received by the buffer, the End-of-Record signal generates an End-of-Record interrupt (if armed) and sets the Inter-record indicator; the channel does not disconnect.

2.101 When the peripheral device is magnetic tape, the tape continues to move after it signals End-of-Record. As in reading tape, the signal causes the Tape Gap signal to come high. If the program executes a new write tape or erase tape EOM during the inter-gap time (approximately one millisecond), the tape remains in motion and proceeds to write or erase a new record. If the program executes no such EOM before the Tape Gap signal drops, the channel disconnects and the tape comes to a stop. No interrupt occurs at this time. This is the only condition which causes a channel to disconnect automatically.

2.102 To proceed after the End-of-Record occurs, the program first executes a Buffer Control mode EOM to re-initialize the Channel Unit Address Register and then reloads the interlace portion of the channel (the program can alert the Interlace via the Buffer Control

EOM). Otherwise, the channel immediately terminates any attempt to use its interlace portion, since the channel is aware that it is still active and in the End-of-Record condition. When the program continues from an Inter-record condition, the program should use an extended mode terminal function.

2.103 A program should not use IORP with devices that do not generate End-of-Record responses upon output termination; such devices are paper tape and typewriter. These devices do terminate output but give the program no indication when they receive the last characters.

2.104 The IORP should also not be used with the printer and card punch since these devices expect the channel to disconnect after they send EOR.

2.105 INPUT/OUTPUT UNTIL SIGNAL THEN PROCEED (IOSP)

2.106 Input

2.107 Read C words. If the channel counts C down to zero before the peripheral device encounters the End-to-Record, the channel generates a Count Equals Zero interrupt (if armed). The program should reload the interlace portion of the channel to continue reading the record. As far as the peripheral device knows, nothing happens at this time. Failure to reload the Interlace before the peripheral device sends enough characters to overflow the channel buffer causes a rate error; this sets the channel error indicator.

2.108 When the peripheral device encounters the End-of-Record, IOSP operates identically like the IORP command.

2.109 Output

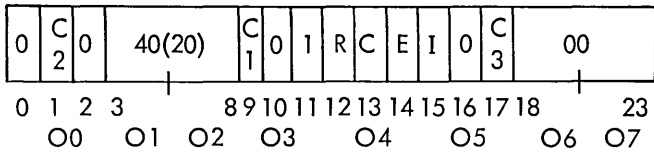
2.110 Write C words. When the channel counts C down to zero, the channel generates a Count Equals Zero interrupt (if armed); the channel does not terminate output. The program should reload the interlace portion of the channel to continue writing in the same record. Failure to reload the Interlace before the buffer transmits all of the characters in its registers and before the peripheral device requests the next character from the buffer results in a rate error; this sets the channel error indicator.

2.111 If the program executes a TERMINATE OUTPUT (TOP) instruction after the channel has counted C down to zero, the channel terminates the output and operates identically like the IORP from this point on.

2.112 CHANNEL AND DEVICE SKS

2.113 The Channel and Device Test mode SKIP IF SIGNAL NOT SET (SKS) tests the indicators in a channel as well as devices attached to it. To test the channel, use unit address 00. The instruction format is:

2.114 CHANNEL TESTS

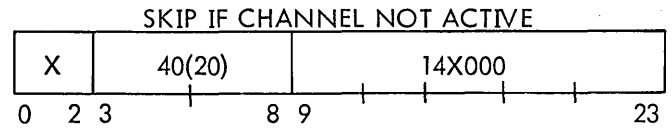


Bit Designation	Octal Position	Octal Value	Function
40(20)	O1-O2	40(20)	Bit positions 3 through 8 contain 40(20), the SKS instruction code.
01	O3	1	Bit positions 10 and 11 contain the mode selection.
C1	O3	4	Bits C1, C2, C3 used as an octal address, specify the channel to be tested. Channel W (A) is 0, channel Y (B) is 1, and so on.
C2	O0	2	
C3	O5	1	
R			Test for ready. A 1-bit selects the test. Skip if Ready or Inactive.
C	O4	2	Test if indicator for Word Count Equal to Zero is set. A 1-bit selects the test. Skip if word count zero.
E		1	Test for error indicator reset. A 1-bit selects the test. Skip if no error.
I	O5	4	Test for Inter-record condition.
00			Bit positions 18 through 23 are zero to specify a channel test. Each of these tests causes a skip when the test condition is true.

2.115 STANDARD SKS INSTRUCTIONS

2.116 Several SKS function configurations have standard uses. These have standard, assembler-type mnemonics and are always used as shown.

2.117 CHANNEL ACTIVE TEST (CAT)



2.118 If the channel is ready to accept a new input/output instruction, the computer skips the next instruction in sequence and executes the following instruction. If the channel is active, or in the process of disconnecting a peripheral unit, the computer executes the next instruction in sequence.

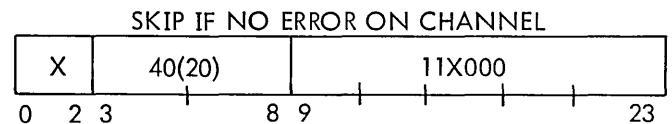
Mnemonic	Channel Active Test	Instruction
CAT 0	W (A)	0 40(20) 14000
CAT 1	Y (B)	0 40(20) 14100
CAT 2	C	2 40(20) 14000
CAT 3	D	2 40(20) 14100

2.119 The following SDS 900 series compatible instructions make the identical test as the above instructions on channels W and Y;

BRTW	0 40 21000	W BUFFER READY TEST
BRTY	0 40 22000	Y BUFFER READY TEST

2.120 The indicator that CAT tests is reset only by the next EOM that connects and alerts the same channel.

2.121 CHANNEL ERROR TEST (CET)



2.122 CET tests the error indicator in the channel for being in the set condition. If the error indicator has not been set, the computer skips the next instruction in sequence and executes the following instruction. If the error indicator has been set, the computer executes the next instruction in sequence.

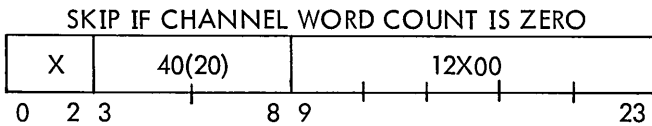
Mnemonic	Channel Error Test	Instruction
CET 0	W (A)	0 40(20) 11000
CET 1	Y (B)	0 40(20) 11100
CET 2	C	2 40(20) 11000
CET 3	D	2 40(20) 11100

2. 123 The following SDS 900 series compatible instructions make the identical test of channels W and Y:

BETY 0 40 20020 Y BUFFER ERROR TEST
 BETW 0 40 20010 W BUFFER ERROR TEST

2. 124 The indicator that CET tests is reset only by the next EOM that connects and alerts the same channel.

2. 125 CHANNEL ZERO COUNT TEST (CZT)

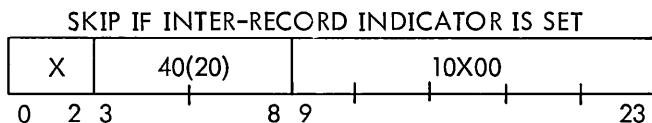


2. 126 CZT tests whether the contents of the Word Count Register in the channel have been reduced to zero. If the contents of WCR are zero, the computer skips the next instruction in sequence and executes the following instruction. If the contents of the WCR are non-zero, the computer executes the next instruction in sequence.

Mnemonic	Channel Zero Count Test	Instruction
CZT 0	W (A)	0 40(20) 12000
CZT 1	Y (B)	0 40(20) 12100
CZT 2	C	2 40(20) 12000
CZT 3	D	2 40(20) 12100

2. 127 The indicator that CZT tests is reset only by a POT instruction to set up the word count and data address in the same channel.

2. 128 CHANNEL INTER-RECORD TEST (CIT)



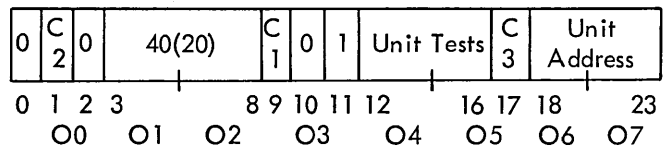
2. 129 CIT tests the Inter-record indicator in the selected channel. If the Inter-record indicator is set, the computer skips the next instruction in sequence and executes the following instruction. If the indicator is reset, the computer executes the next instruction in sequence.

Mnemonic	Channel Active Test	Instruction
CIT 0	W (A)	0 40(20) 10400
CIT 1	Y (B)	0 40(20) 10500
CIT 2	C	2 40(20) 10400
CIT 3	D	2 40(20) 10500

2. 130 The Inter-record indicator is set only during extended mode operation when using a Proceed Function; the indicator is set for an inter-record or zero count condition. The indicator is reset by the next alert and connect EOM.

2. 131 DEVICE TESTS

2. 132 The SKIP IF SIGNAL NOT SET (SKS) below, used in the Channel and Device Test mode, tests the condition of the peripheral devices in the system directly. The peripheral device sections contain the individual instruction descriptions.



Bit Designation	Octal Position	Octal Value	Function
C1	O3	4	Bit positions 9, 1, and 17 are used as an octal digit to specify the channel.
C2	O0	2	Channel W (A) is 0, channel Y (B) is 1, and so on.
C3	O5	1	
40(20)	O1-O2	40(20)	Bit positions 3 through 8 contain the SKS instruction code.
01	O3	1	Bit positions 10 and 11 contain the mode selection.
Unit Tests	O4-O5		Bit positions 12 through 16 select the particular test and are system dependent.
Unit Address	O6-O7		Bit positions 18 through 23 specify the unit address.

location processes the interrupt. This routine contains the instruction (MIW or WIM, for example) that can execute immediately without computer tie-up.

2. 153 During single-word operations, a parity error or incorrect timing error sets the buffer error indication in the channel. The incorrect timing error occurs when characters enter the buffer during input before the removal of the previous word; during output, buffer error indication occurs if characters are needed for output before the buffer receives the next word. The transmission does not terminate upon detection of any of these errors.

2. 154 The interrupt system can detect an End-of-Record termination. During output, use of TERMINATE OUTPUT (TOP) after the final MIW (MIY) causes an interrupt to the appropriate End-of-Transmission location when that final data word has been processed by the buffer. This interrupt takes the place of the End-of-Word interrupt; the End-of-Transmission condition inhibits the End-of-Word interrupt. During input, the End-of-Transmission interrupt is sent to the End-of-Transmission location when the End-of-Record is detected. During input from devices which do not generate an End-of-Record, an EOM disconnects (DSC) the channel to terminate the transmission. This termination generates no End-of-Transmission interrupt.

SECTION III THEORY OF OPERATION

3.1 GENERAL

3.2 The TMCC communicates with an external system or device by means of a shift register utilizing 6, 12, or 24-bit characters plus a parity bit. The maximum character size depends on the optional registers that may be installed. In TMCC Models 922XX, the character length is fixed at 6, 12, or 24 bits and can not be varied from one size to another. However, TMCC Models 932XX having the larger character length options may be switched from one size to another under control of the external system. The length is selected through activation of one of the character size control lines. External devices may activate these lines as necessary but if none are activated the TMCC assumes 6-bit characters. The rate of information transfer is determined by a clock signal from the external device. For both input and output, the TMCC slaves itself to the clock frequency of the device up to the TMM maximum data rate of two machine cycles per character.

3.3 The TMCC communicates with the computer by means of a 24-bit shift register which transfers words, an octal group at a time, between the TMCC and the computer C Register. The TMCC thus has two registers for data storage. These provide the means to assemble input characters into words or disassemble words into output characters. The number of characters per word is under program control but is limited to a maximum of four 6-bit characters, two 12-bit characters, or one 24-bit character.

3.4 Information may be input or output by executing an instruction for each word (channels W (A), or Y, only). The instruction may be given in advance of the time it is needed, in which case the computer remains idle until the channel is ready. Or the computer interrupt system may be employed so that the channel can call for an instruction when it is ready to use it. This allows the computer to continue with other computations when not actually engaged in the input/output (I/O) process.

3.5 An optional interlace feature may be installed in the TMCC to facilitate I/O operations with fewer instructions. The interlace logic allows a program to designate to the TMCC how many words are to be transferred and the memory location of the first word. Then, without further instructions, the TMCC can assemble or disassemble the number of words specified and time-share with the computer each time it is ready to transfer a word to or from memory. The I/O process may thus be interlaced with computation or with similar I/O operations on other channels.

3.6 Figure 3-1 is a block diagram of the TMCC. As indicated on the figure, the principal parts of the TMCC are the Character Buffer, Word Assembly Register, Unit Address Register, Word Counter, and Address Counter.

3.7 CHARACTER BUFFER

3.8 Depending on the option, the Character Buffer is a single character storage register of 6, 12, or 24 bits. It is implemented with S-R type flip-flops connected in a series-parallel manner so shifting takes place in octal groups (i. e. three bits in parallel). Transfer of data between the Character Buffer and an external system is entirely parallel for the whole character whether it be input or output.

3.9 WORD ASSEMBLY REGISTER

3.10 This register is also connected in series-parallel to allow shifting in octal groups. It is composed of three parallel registers of eleven flip-flops each. The register includes twenty-four flip-flops to store a complete word plus additional flip-flops on the ends of each series string to allow for timing considerations. The Word Assembly Register is implemented with flip-flops connected in such a manner that continuous recirculation occurs.

3.11 UNIT ADDRESS REGISTER

3.12 The 6-bit address code to select a specific peripheral device is set-up in the Unit Address Register. The register is composed of five flip-flops whose outputs are sent to the peripheral unit for decoding. A sixth flip-flop, associated with the Unit Address Register, is also set-up at the same time to signal the external device whether an input or an output is to take place.

3.13 WORD COUNTER

3.14 The Word Counter is part of the optional interlace equipment. It is a fifteen stage flip-flop counter used to store the number of words to be transferred during an interlaced I/O operation. With each word transfer, the counter is decremented by "one". (Actually, the complement of the count is incremented.)

3.15 ADDRESS COUNTER

3.16 This counter is also part of the optional interlace feature. Its purpose is to store the address of the d memory location currently being accessed. Each time a word is taken from or sent to memory, the Address

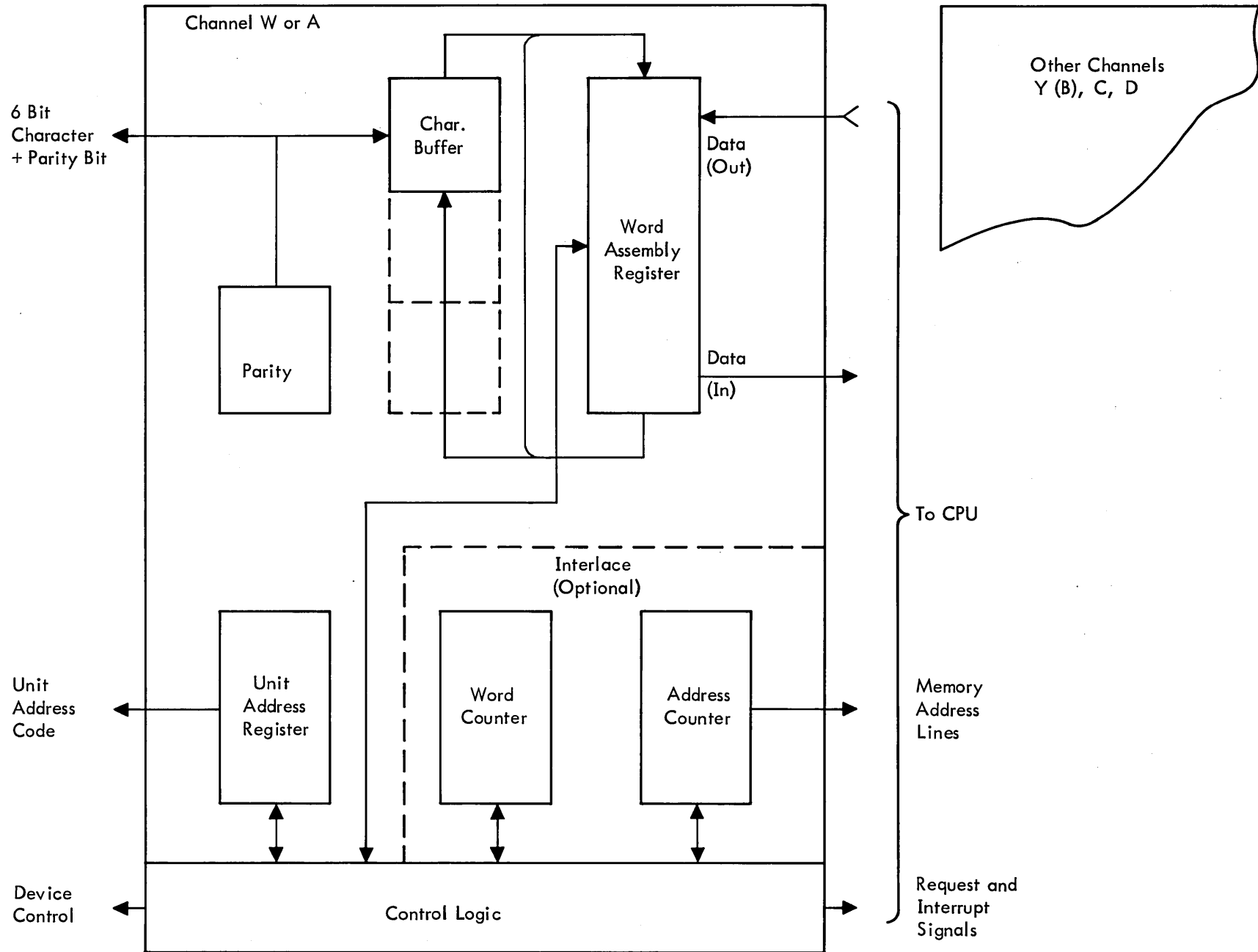


Figure 3-1. Time Multiplexed Communication Channel With Interlace, Block Diagram

Counter is incremented by "one" to prepare for locating the next word.

3. 17 INPUT PROCESS

3. 18 A brief outline of a non-interlaced input process follows. The TMCC is initialized by an EOM instruction to set-up the address code of the peripheral device, to designate input or output, and specify the number of characters per word. When the address is decoded by a peripheral device, the device is activated and begins sending clock signals and data to the channel.

3. 19 Input clocks are synchronized with TMCC timing by clock counter flip-flops which detect the peripheral device clock, transfer the input data to the Character Buffer and cause the Word Assembly Register to circulate through the Character Buffer. Ordinarily the Word Assembly Register recirculates on itself but as each new input character is ready it circulates through the Character Buffer for one machine cycle to pick up the new character. A machine cycle is defined as eleven computer clock pulse periods.

3. 20 As each character is clocked in, the character counter (W7 W8) is decremented until the count reaches zero, signaling that a complete word has been formed. Depending on the particular channel (TMCC W (A) or Y) a WIM, AIM, or YIM instruction may be used to transfer the word from the TMCC to memory. When the instruction is executed, a word is shifted in octal groups from the Word Assembly Register to the computer C Register and then stored in the memory location specified by the effective address of the WIM, AIM, or YIM instruction.

3. 21 To prepare for the next word, the Character Counter is reset to its original count which was designated by the initializing EOM instruction. Storage of the original count is accomplished by utilizing excess positions of the Word Assembly Register.

3. 22 The input process may be terminated by another EOM instruction or by detection of an externally applied halt condition.

3. 23 OUTPUT PROCESS

3. 24 A non-interlaced output process is started in the same way that an input operation is initialized. A MIW (MIY or MIA) instruction can then be used to transfer words from memory to the Word Assembly Register via the computer C Register.

3. 25 Each time the external device sends a clock signal to the TMCC, it is detected by the Clock Counter.

At a certain point in the counter sequence, the device extracts information from the Character Buffer. At the end of the counter sequence, the Word Assembly Register recirculates through the buffer for one machine cycle. At the conclusion of the cycle, a new character is available in the Character Buffer awaiting the next clock. The Clock Counter also decrements the Character Counter. When it is decremented to zero, the last character of a word is in the buffer. By this time another MIW (MIY or MIA) instruction should be executed; or the instruction may be called for through the interrupt system. As with input, the Character Counter is reloaded between words from the Word Assembly Register.

3. 26 The output operation is concluded by a Terminate Output EOM instruction.

3. 27 Parity

3. 28 During input, a parity flip-flop accepts the input parity bit and checks the character for odd parity. The same flip-flop is also used to generate the parity bit for output characters.

3. 29 TMCC AND INTERLACE CONFIGURATION

3. 30 The theory of operation contained in this manual is applicable to all models of the Time-Multiplexed Communication Channels for the 925/930/9300 Computers. The operation is for the most part identical for all models. Those differences that do exist, however, are covered in the explanation where appropriate.

3. 31 DETAILED DESCRIPTION

3. 32 Many portions of the TMCC logic are common to both input and output operations. Other portions of the logic are specifically related to input only, output only, or interlace only. The following logic description begins with the common logic functions.

3. 33 Subsequent to the logic descriptions are the Glossary of Logic Term and the Logic Equations.

3. 34 PULSE COUNTER

3. 35 The pulse counter consists of flip-flops Qr1, Qr2, Qr3, and Qr4. These flip-flops perform the same function for the TMCC that the pulse counter, Q1 through Q6, does for the central processor (CPU). The counter is included in the TMCC in addition to the counter in the CPU in order to avoid excessive delays and loading that would occur if all timing signals were obtained directly from the CPU. The Pulse Counter defines the pulse times, T8, T7, T6, T5, T4, T3, T2, T1, T0, Tr, and Tp needed for timing of all processes within the unit. Refer to table 3-1 for the Pulse Counter Truth Table.

Table 3-1. Pulse Counter Truth Table

	Qr1	Qr2	Qr3	Qr4
Tp	0	0	0	0
T8	1	0	0	0
T7	1	0	0	1
T6	1	1	1	1
T5	0	1	1	1
T4	0	0	1	1
T3	0	0	1	0
T2	1	0	1	0
T1	1	1	1	0
T0	0	1	1	0
Tr	0	0	0	0
Tp	0	0	0	0

3.36 It is logically impossible to set Qr2, Qr3, and Qr4 unless Qr1 has been previously set. Therefore, without the additional set signal provided by Tpc, the counter would soon reach the reset state of time Tr + Tp and then stop counting. However, Tpc, the Tp signal from the CPU, is used to advance the TMCC pulse counter from Tp to T8. Thus, both counters advance to T8 at the same time and both remain in synchronism. An examination of the logic also indicates that the two pulse counters will synchronize regardless of the turn-on state.

$$\begin{aligned} sQr1 &= Tpc + \overline{Qr2} \overline{Qr3} \overline{Qr4} \\ rQr1 &= Qr2 \\ sQr2 &= Qr1 \overline{Qr2} \overline{Qr4} (Qr4 + T0) + Qr1 \overline{Qr3} \\ rQr2 &= \overline{Qr1} \\ sQr3 &= Qr1 \overline{Qr4} \\ rQr3 &= Qr3 \overline{Qr4} (Qr4 + T0) \\ sQr4 &= Qr1 \overline{Qr3} \\ rQr4 &= \overline{Qr1} \overline{Qr2} \end{aligned}$$

Only those pulse times or combinations of pulse times that are need for timing in the TMCC are decoded. For example:

$$\begin{aligned} T0 &= \overline{Qr1} \overline{Qr2} \overline{Qr4} \\ T6 - T0 &= Qr3 \end{aligned}$$

3.37 In examining the pulse counter and the T0 decoding logic, it appears that the term (Qr4 + T0), as used on the inputs of Qr2 and Qr3, contains redundant logic. This is a result of logic mechanization and the redundant terms are not significant to the counter operation.

3.38 A TMCC pulse counter is associated with each W and C channel. The Y and D channels share the same pulse counters as the W and C channels, respectively. Other logic is shared in a similar manner by two channels and will be noted as each case arises.

3.39 INPUT/OUTPUT PROCESSING

3.40 When an EOM instruction of the form, EOM-0XXXX or EOM4XXXX, is executed to start an input or output process, a Buffer Control Signal, Buc, produces a clear signal, Wc, and a set signal, Ws. These two signals, Wc and Ws, permit initialization for the input/output operation. The Unit Address and the Character Count are set up from the C-Register. The registers are first cleared by Wc and then set by Ws. Refer to figure 3-2.

$$\begin{aligned} Buc &= Eom \overline{C10} \overline{C11} \overline{C1} * \\ Wc &= Buc \overline{C17} (T6 + T5) + St + \dots \\ Ws &= Buc \overline{C17} (T3 - T0) \end{aligned}$$

*C1 is used in place of $\overline{C1}$ for Channels C and D. Similarly, C17 and $\overline{C17}$ distinguish between the W and Y channels, or between the C and D channels.

(T6 + T5) and (T3 - T0) are decoded timing signals from the pulse counter. Frequent use is made of such pulse times throughout the manual without further explanation. The term $\overline{C17}$ distinguishes enabling of the W channel rather than the Y channel. The combination of terms which make up the input of Buc indicate that the computer is in phase (Ø) 5 of the execution of an EOM Buffer Control instruction. The term, St, is produced by the start switch which may also be used to reset the TMCC. The register logic involved in the clear and set sequencing is:

$$\begin{aligned} sW14 &= Ws C23 \\ rW14 &= Wc \\ sW11 &= Ws C20 \\ rW11 &= Wc \\ sW10 &= Ws C19 + \dots \\ rW10 &= Wc \end{aligned}$$

Unit Address Register
5-bits provide 31 addressing codes and a disconnect code.

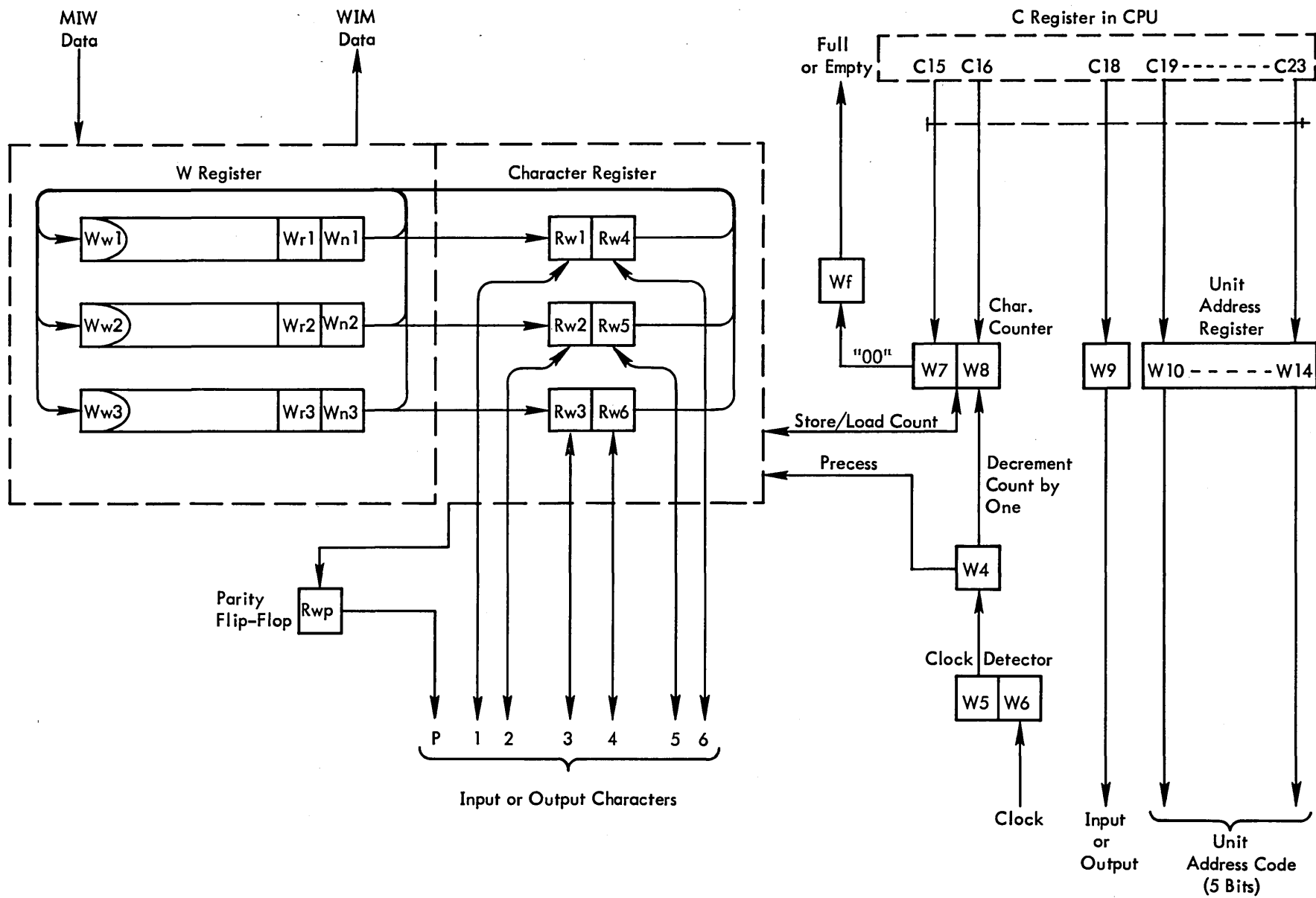


Figure 3-2. TMCC Information Flow Diagram, Input/Output (6 bit)

$$sW9 = Ws C18 \quad W9 = 1 \text{ for output}$$

$$rW9 = Wc \quad W9 = 0 \text{ for input}$$

$sW8 = Ws C16 + \dots$	} <u>Character Counter</u>		W7	W8	Characters/Word
$rW8 = Wc + \dots$			1	1	4
$sW7 = Ws C15 + \dots$	1	0	3		
$rW7 = Wc + \dots$	0	1	2		
	0	0	1		

3.41 At the end of the Buc type instruction and during the last pulse time (T0) that Ws is on, the flip-flop, W4, is set on for three pulse times.

$$sW4 = Ws T0 + \dots$$

$$rW4 = W4 T8 + \dots$$

During this time, W4 allows the contents of the character counter to be copied into the W-Register. The bit in W7 goes into Ww1 and the bit in W8 goes into Ww2.

$$Ww1 = W4 W7 (\overline{T7 - T0}) + \dots$$

$$Ww2 = W4 W8 (\overline{T7 - T0}) + \dots$$

During this period, $\overline{W4}$ or timing signals inhibit the other inputs to the W-Register.

3.42 This process provides a means for the TMCC to remember the initial character count while using flip-flops W7 and W8 to perform the actual count-down of the characters as each word is assembled on input or disassembled on output. As a word is processed, the original count is reloaded into the counter to get ready for the next word.

$$sW8 = Wxx Wn2 (\overline{T7 - T0}) \overline{W4}$$

$$sW7 = Wxx Wn1 (\overline{T7 - T0}) \overline{W4}$$

The term Wxx is true during the transmission of a word being processed. Thus, in addition to its normal word handling functions, the W-Register provides storage of the character count for W7 and W8.

3.43 W-Register

3.44 The W-Register is a one word recirculating flip-flop register with additional flip-flops at the read and write ends. The Write input signals to the first flip-flops of the register are designated Ww-, the intermediate Read flip-flops are designated Wr- and these drive the last flip-flops which are designated Wn-. The n represents now.

3.45 During normal recirculation, the outputs of the n flip-flops are fed to the w inputs, which, in turn, feed the remaining serial shift circuits. There are a total of eleven stages in a recirculating loop, one for each pulse time of a machine cycle. Three recirculating loops are required to hold a 24-bit word. The first recirculating loop holds only the most significant bits of each of the eight octal digits in a word. The second recirculating loop holds the middle bits of the eight octals. The third recirculating loop holds the least significant bits of each of the eight octals. If the octal word 07030407 (000 111 000 011 000 100 000 111) were being held in the register, its bits at pulse time Tp would appear as follows:



At the next pulse time, T8, all bits would be moved one bit position to the right and for succeeding pulse times as the register recirculates. At the least significant pulse time, T7, the bits of the least significant octal digit are in their respective now flip-flops, Wn1, Wn2, and Wn3. At succeeding pulse times the now flip-flops present the corresponding octal digits. The additional positions beyond the eight octals of the register act as fill-in bits to satisfy logic timing requirements (so that a word is back in its original position after one machine cycle of 11 clock pulses) and are also used to store the character count. Normal recirculation is allowed by the following logic:

$$Ww1 = \overline{W4} Wn1 \overline{Wxx} + \dots \quad \overline{W4} \overline{Wxx} \text{ indicates the W-Register is not participating in some other operation.}$$

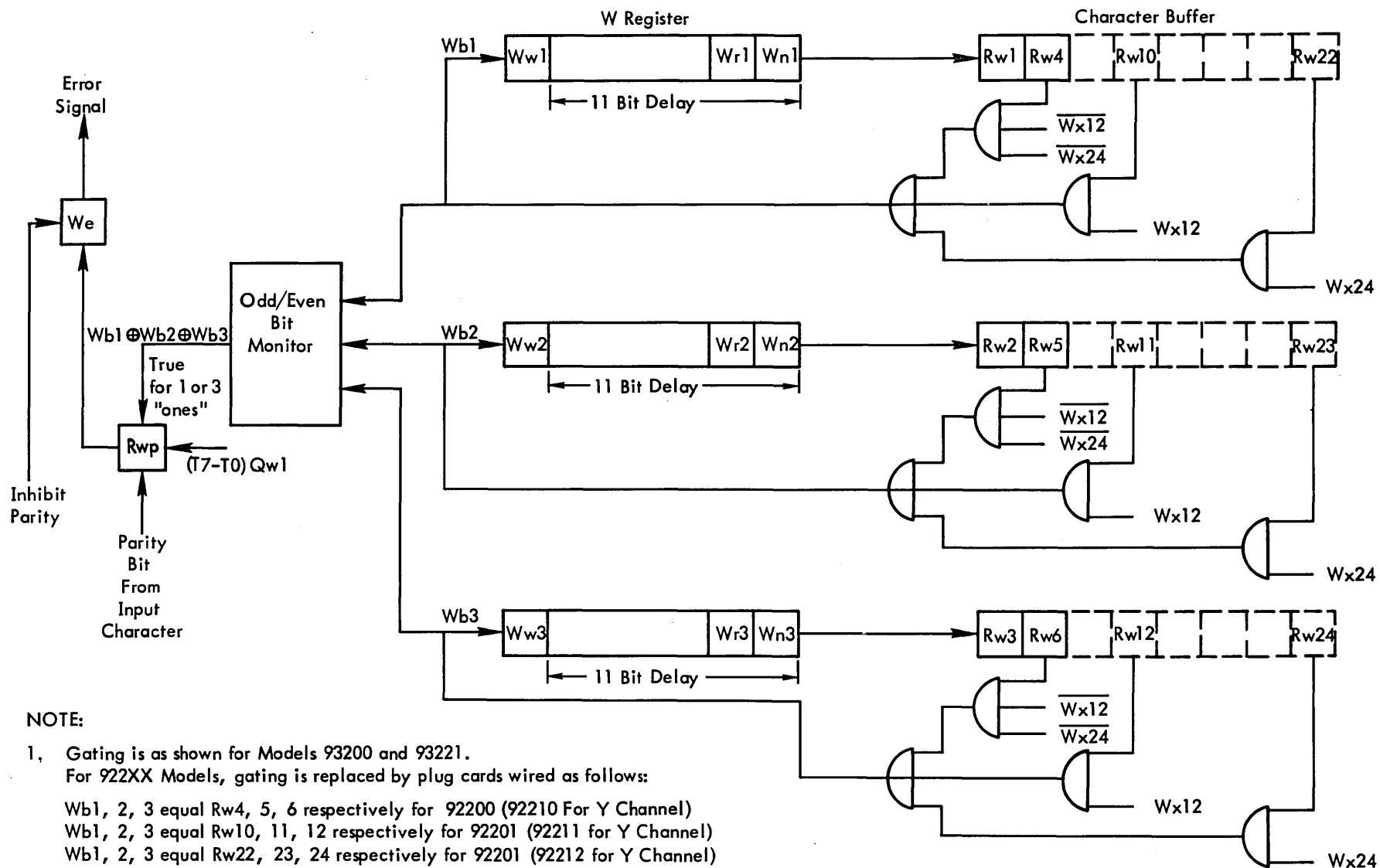
$$Ww2 = \overline{W4} Wn2 \overline{Wxx} + \dots$$

$$Ww3 = \overline{W4} Wn3 \overline{Wxx} + \dots$$

3.46 The Character Register is composed of six, twelve, or twenty-four R-S type flip-flops. When shifting an input character from the Character Register into the W-Register or an output character from the W-Register to the Character Register, W4 is set for the period (T7 - T0) of one machine cycle. This enables the gating of one character precession in the W-Register by causing the data in the W-Register to recirculate through the Character Register. Refer to figure 3-3.

$$Ww1 = W4 Wb1 (T7 - T0) + \dots$$

$$Ww2 = W4 Wb2 (T7 - T0) + \dots$$



NOTE:

1. Gating is as shown for Models 93200 and 93221.
For 922XX Models, gating is replaced by plug cards wired as follows:
Wb1, 2, 3 equal Rw4, 5, 6 respectively for 92200 (92210 For Y Channel)
Wb1, 2, 3 equal Rw10, 11, 12 respectively for 92201 (92211 for Y Channel)
Wb1, 2, 3 equal Rw22, 23, 24 respectively for 92201 (92212 for Y Channel)
2. Wx12 and Wx24 are signals derived from the peripheral device to select the I/O character size. If the device does not designate a character size then $\overline{Wx12}$ and $\overline{Wx24}$ are true and Wx12 and Wx24 are false.

Figure 3-3. Precession Loop and Input Parity Checking Logic

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$$\begin{aligned}
 Ww3 &= W4 Wb3 + \dots \\
 sRw1 &= W4 \overline{Wxx} (T7 - T0) Wn1 + \dots \\
 rRw1 &= W4 \overline{Wxx} \overline{Wn1} + \dots \\
 sRw2 &= W4 \overline{Wxx} (T7 - T0) Wn2 + \dots \\
 rRw2 &= W4 \overline{Wxx} \overline{Wn2} + \dots \\
 sRw3 &= W4 \overline{Wxx} (T7 - T0) Wn3 + \dots \\
 rRw3 &= W4 \overline{Wxx} \overline{Wn3} + \dots
 \end{aligned}$$

$$\begin{aligned}
 sRw24 &= W4 \overline{Rw24} R w21 + \dots \\
 rRw24 &= W4 R w24 \overline{Rw21} + \dots
 \end{aligned}$$

The shift inputs for $Rw1$, $Rw2$, and $Rw3$ are given in paragraph 3. 46. The precession of characters is controlled by $W4$. When a WIM or MIW instruction is executed, Wf is set. Wf remains set until the count in $W7$ and $W8$ is 00.

$$\begin{aligned}
 sWf &= Rx T0 Twy \\
 rWf &= \overline{W7} \overline{W8} W4 (T6 + T5)
 \end{aligned}$$

$Rx Pwy = Wxx$ which indicates a WIM or MIW instruction is occurring. Wf is later used to indicate to the computer when the W -Register is full on input or empty on output.

3.50 Two other flip-flops, $W5$ and $W6$, detect external device clocks (which may occur either before or after Wf is set). The clock Ecw is first detected by $W6$ as follows:

$$\begin{aligned}
 sW6 &= \overline{W5} Ecw T8 \overline{W10} \overline{W11} \overline{W12} \overline{W13} \overline{W14} \\
 rW6 &= W4 T0 + Wc
 \end{aligned}$$

At the next $T0$ pulse after the device clock goes false, $W5$ sets allowing $W6$ to reset again. Setting of $W5$ inhibits further clock detection until the current clock is processed.

$$\begin{aligned}
 sW5 &= \overline{W5} W6 \overline{Ecw} T0 + \dots \\
 rW5 &= W4 T0 + \dots
 \end{aligned}$$

Setting of $W4$ is interlocked with $W5$ and Wf to allow one precession of the W Register to take place after each external clock signal.

$$\begin{aligned}
 sW4 &= W5 Wf T8 \overline{Wg} + \dots \\
 rW4 &= W4 T0 + \dots
 \end{aligned}$$

The term, \overline{Wg} , explained in detail in paragraphs 3-66 and 3-150) is assumed to be true at this time.

3.51 The equation for Wf contains $\overline{W7} \overline{W8}$, indicating that precession is complete when the character count reaches 00. The count down is enabled by $W4$ and is accomplished by:

$$\begin{aligned}
 sW8 &= W7 \overline{W8} W4 T0 + \dots \\
 rW8 &= W8 W4 T0 + \dots \\
 rW7 &= W7 \overline{W8} W4 T0 + \dots
 \end{aligned}$$

3.52 The foregoing discussion was limited to processes common to input and output operations. The interrupt

Wxx is a signal denoting that an MIW, WIM or Time Share operation is occurring.

$$\begin{aligned}
 Wxx &= Rx Pwy + \dots * \\
 Pwy &= 05 **
 \end{aligned}$$

* Rx is always false for the C-channel. Another term, not shown allows Wxx to function in that channel only on an interlaced basis.

**In Model 92200. See 05, $Pw5$ in the Glossary of terms.

3.47 The Time Share Operation is discussed beginning with paragraph 3. 102. Since there are no WIM/MIW type instructions for TMCC-B, C, or D, Rx is always false for these channels.

3.48 The timing signals $(T7 - T0)$ on the inputs of $Ww1$ and $Ww2$ are necessary to prevent interference with the storage of the character count which was previously described. The character count does not precess but recirculates. This recirculation is allowed by:

$$\begin{aligned}
 Ww1 &= \overline{W4} Wn1 (\overline{T7 - T0}) + \dots \\
 Ww2 &= \overline{W4} Wn2 (\overline{T7 - T0}) + \dots
 \end{aligned}$$

3.49 During the input or output precession, the shift logic for the character register is as shown below. Depending on the 12- or 24-bit character option, all bits may not be installed.

$$\begin{aligned}
 sRw4 &= W4 \overline{Rw4} R w1 + \dots \\
 rRw4 &= W4 R w4 \overline{Rw1} + \dots \\
 sRw5 &= W4 \overline{Rw5} R w2 + \dots \\
 rRw5 &= W4 R w5 \overline{Rw2} + \dots \\
 &\quad \vdots \qquad \qquad \quad \vdots \\
 sRw23 &= W4 \overline{Rw23} R w20 + \dots \\
 rRw23 &= W4 R w23 \overline{Rw20} + \dots
 \end{aligned}$$

function, which allows the TMCC to signal the computer when it is ready to transfer a word to or from memory, was omitted. In subsequent paragraphs, features peculiar to input or output operation are discussed separately. Details concerning interrupt operation are included.

3. 53 INPUT PROCESS (W9 true)

3. 54 The characteristics of a typical input clock signal and its relationship to the input data and Clock Counter flip-flops are illustrated in figure 3-4. Examination of figure 3-4 indicates that if the data is to be sampled by W6 W5 it has to be on by the time the clock, Ecw, returns to zero volts. Also, in order not to be read by the previous or next clock, the data may not come on until at least one machine cycle after the previous clock and must go off before the next clock appears.

3. 55 Two detailed examples of the clock signal, Ecw, are illustrated in figure 3-5. To be assured of a timing pulse occurring during the on period (to set W6) and the off period (to set W5) of Ecw, each of these periods must be at least one machine cycle in length. This prescribes an input clock cycle of no less than two machine cycles. The clock rate must be somewhat slower than this for proper operation of W6, W5, and W4. Figure 3-5a illustrates two input clocks with timing such that the second clock is missed. Any clock rate slower than that illustrated in figure 3-5a would be satisfactory, however, a safety margin must be provided to compensate for noise and variations in waveshape and frequency.

3. 56 The input frequency may be increased as illustrated in figure 3-5b if the clocks are interlocked within the peripheral device with W5 and W6 from the TMCC such that

$$\textcircled{E}cw = \overline{(\text{Device Clock})} \overline{W5} \overline{W6}$$

In this case, only two machine cycles per input clock are needed. Some of the peripheral device couplers include this interlocking feature.

3. 57 When a Buc type EOM instruction is executed to activate the TMCC for an input operation, the halt detector, Wh, is reset and the computer interlock flip-flop, Wf, is set by the clearing signal, Wc.

$$\begin{aligned} Wc &= Buc \overline{C17} (T6 + T5) + \dots \\ sWf &= Wc \overline{Wh} + \dots \\ rWh &= Wc (T6 + T5) + \dots \end{aligned}$$

This prepares the Character Register to precess the first input character into the W- Register enabling W4 which gates Ww1, Ww2 and Ww3.

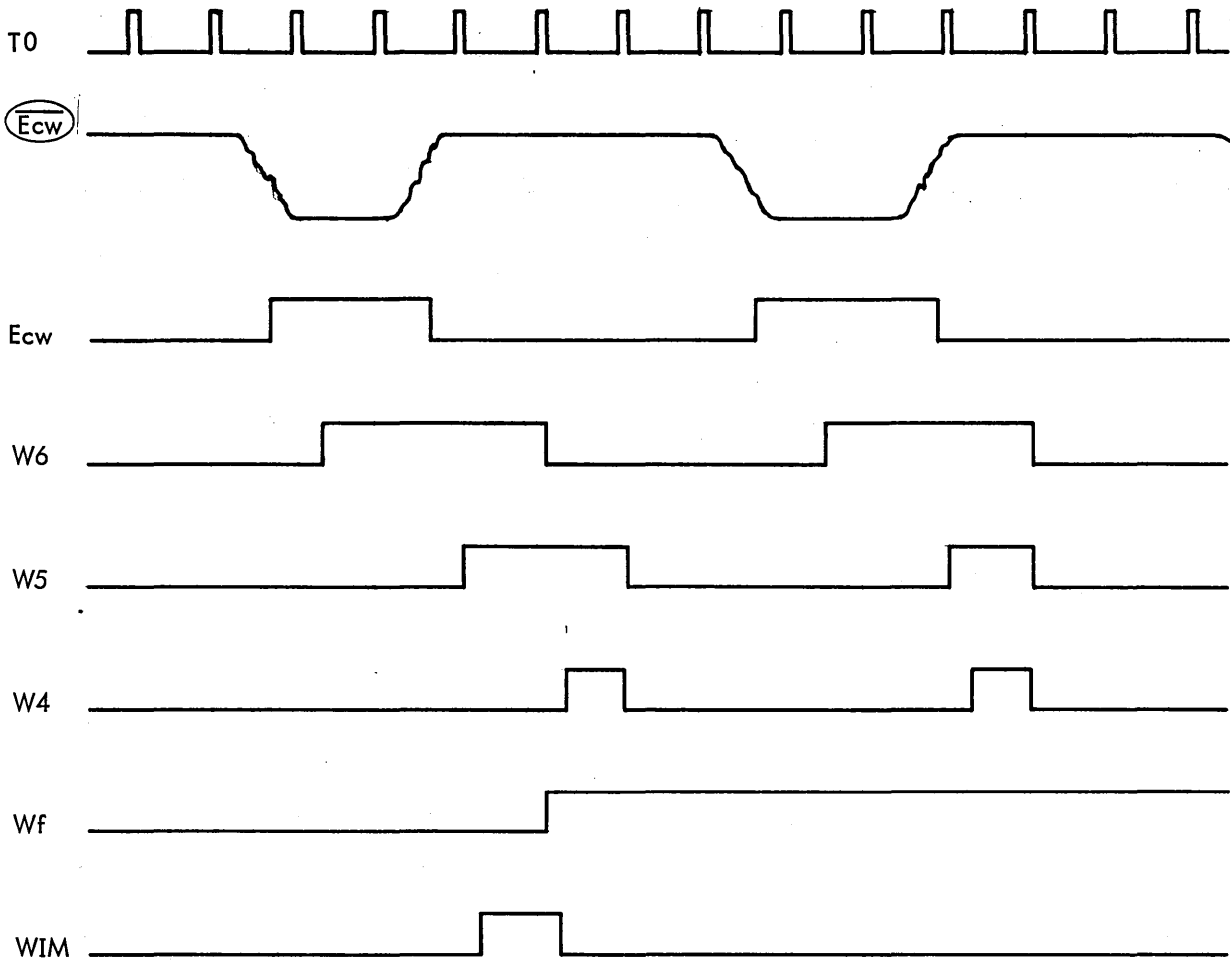
3. 58 The Character Register and parity flip-flop are cleared originally and again between each input character by $\overline{W9} \overline{W6} \overline{W5} \overline{W4}$. After processing the Buc type EOM instruction, the Character Register is ready to receive an input character and clock even though a WIM instruction is not immediately given. The input bits, Zw1 through Zw24 and the parity bit ZwP, are gated into the Character Register and the Parity Flip-Flop by $\overline{W9} \overline{W6} \overline{W5}$. W9 signifies an input operation and W6 W5 indicate an input clock has been detected from the peripheral device.

$$\begin{aligned} sRw1 &= \overline{W9} \overline{W6} \overline{W5} Zw1 \\ rRw1 &= \overline{W9} \overline{W6} \overline{W5} \overline{W4} \\ &\vdots \\ sRw24 &= \overline{W9} \overline{W6} \overline{W5} Zw24 \\ rRw24 &= \overline{W9} \overline{W6} \overline{W5} \overline{W4} \\ \\ sRwp &= \overline{W9} \overline{W6} \overline{W5} ZwP \\ rRwp &= \overline{W9} \overline{W6} \overline{W5} \overline{W4} \end{aligned}$$

3. 59 Precession then takes place as described earlier and another input character can be read into the Character Register. This read-precess cycle is repeated for each input character until the Character Counter is decremented to 00 causing Wf to reset. Resetting of Wf inhibits further precessions and indicates to the computer that the W-Register is full and ready to transfer a word into memory. The program should now provide a WIM instruction to enable the transfer. One more character may still be stored in the Character Register before receiving the WIM instruction. Further input characters are blocked because W4 is inhibited while Wf is reset. This prevents W5 from resetting, which in turn disables clock detection by W6. When a WIM instruction does occur late, The Error Detection Flip-Flop, We, is set. This condition can be tested by an SKS instruction. The error detection flip-flop is reset by Wc which occurs with an EOM-Buc instruction.

$$\begin{aligned} sWe &= W0 \overline{W6} W5 Ecw T8 + \dots \\ rWe &= Wc \overline{Wh} \end{aligned}$$

This equation signifies that a WIM instruction is late if the clock signal, Ecw, is received before the previous character has been precessed out of the Character Register. The W0, Halt Interlock, and Wh, Halt Detector, signals appearing here are discussed in detail in paragraphs 3. 65 and 3. 66. When the interlock feature of figure 3-5 is employed, the condition $\overline{W6} W5 Ecw$ cannot occur to set We. Under such circumstances, the peripheral device (or coupler) must be capable of detecting its own rate errors and reporting them via the Error Signal line, \textcircled{Wes} .



Time →

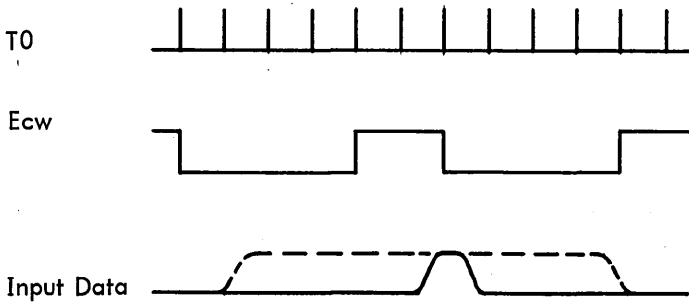
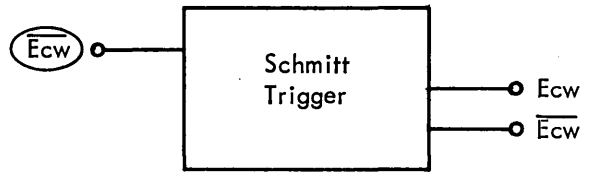
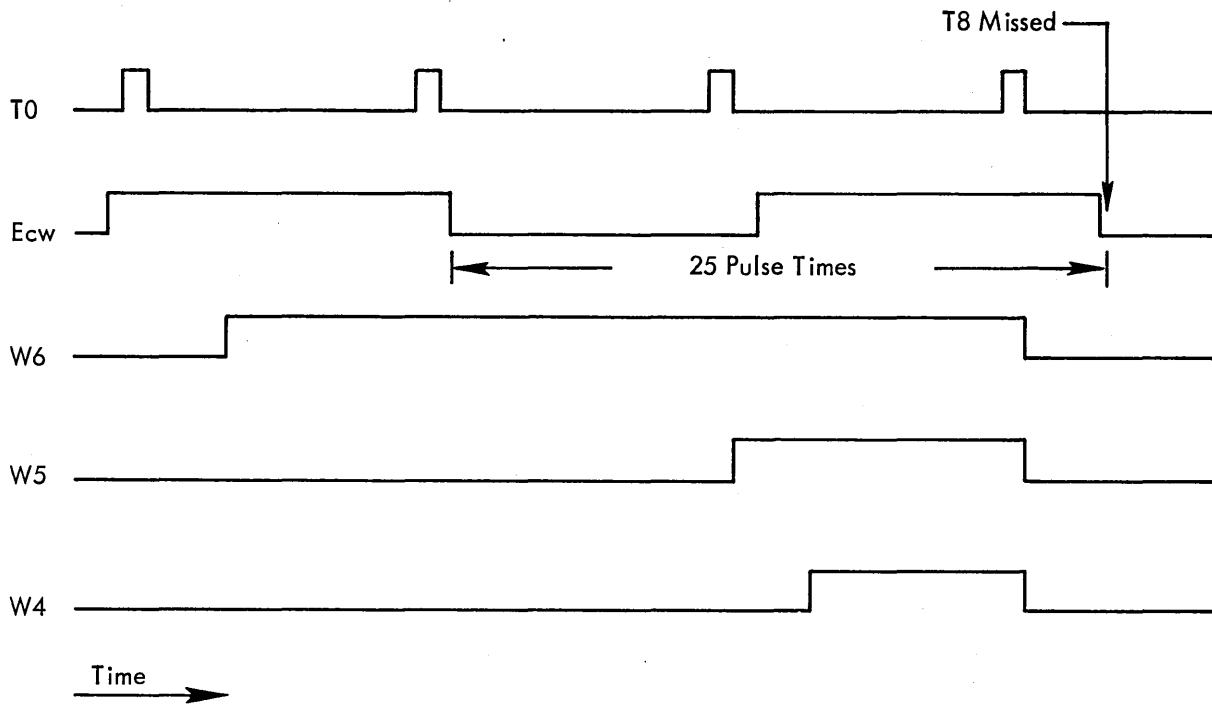
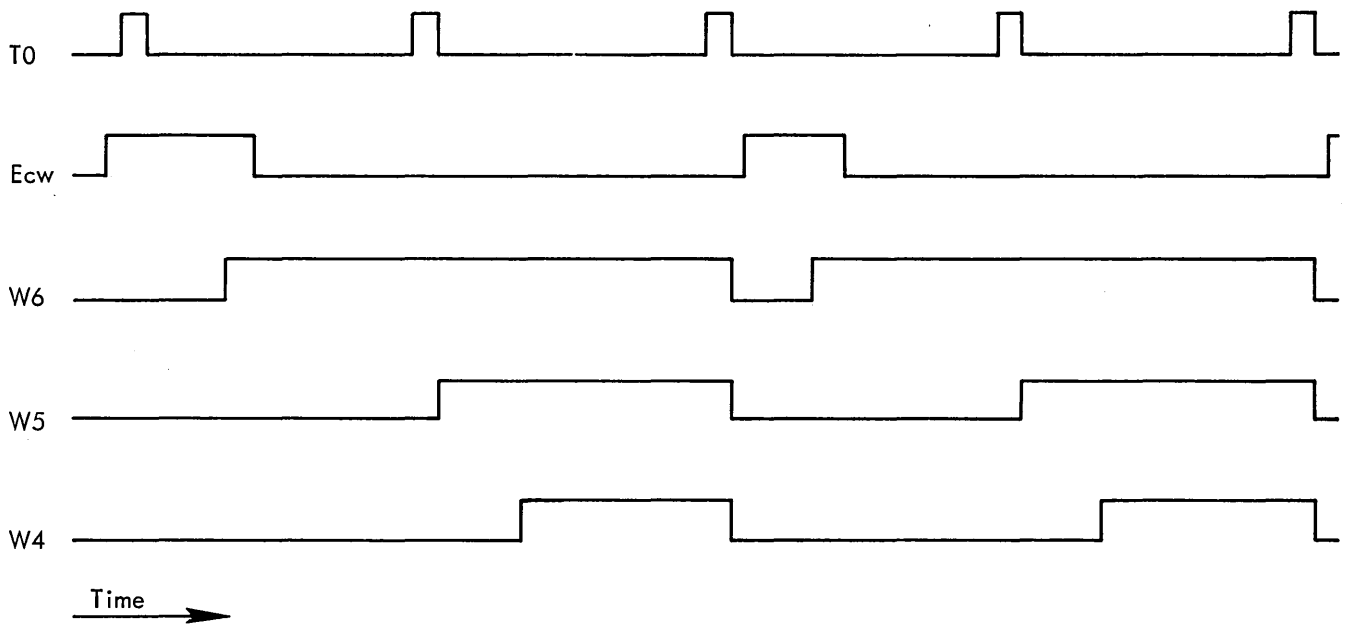


Figure 3-4. Input Clock Timing Charts (Typical)



a. Input Clocks too Fast



b. Input Clocks Interlocked With $\overline{W5} \overline{W6}$

Figure 3-5. Input Clock Timing Charts

$$sWe = Wes + . . .$$

3.60 The Input Timing Chart, figure 3-6, indicates the flow of the basic input process. The first WIM instruction is shown occurring late to illustrate how the Character Register accepts one more character after the W-Register is full.

3.61 The error detector flip-flop may also be set if the input character has even parity.

$$sWe = \overline{W9} \overline{W6} \overline{W5} \overline{W4} Rwp \overline{Wg} Npw \overline{Iwg} + . . .$$

In this equation, $\overline{W9}$ indicates an input process, $\overline{W6} \overline{W5} \overline{W4}$ indicates that the received character has already been precessed into the W-Register when We is set. The term Npw is a signal received from the peripheral device to disable parity checking when a parity bit is not supplied. Wg and Iwg are End-of-Record and Extended Mode signals. These are discussed in paragraph 3-131 dealing with Interlace. Rwp is the Parity Flip-Flop which is toggled while the character is precessing out of the Character Register. During input, the operation of Rwp takes place as follows: Initially and again between each character precession, Rwp is reset by $\overline{W9} \overline{W6} \overline{W5} \overline{W4}$. If the parity bit of the input character is a one, Rwp is set by $\overline{W9} \overline{W6} \overline{W5}$.

$$sRwp = \overline{W9} \overline{W6} \overline{W5} Zwp + . . .$$

$$rRwp = \overline{W9} \overline{W6} \overline{W5} \overline{W4} + Wc + . . .$$

Since W6 is set at pulse time T8 in the above equations Rwp may be set (if Zwp = 1) at the next pulse time, T7. At the next T0, W5 is set and the state of Rwp is reversed at T8, three pulse times later.

$$sRwp = Wf W5 T8 \overline{Rwp} + . . .$$

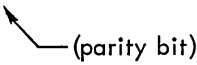
$$rRwp = Wf W5 T8 \overline{W9} Rwp + . . .$$

As a result of this reversal, Rwp is set if the input parity bit was a zero and reset if the bit was a one. The parity of the 6-, 12- or 24-data bits of the input character is now examined. This takes place while the character is being precessed into the W-Register. While precession occurs, the three bits of each octal group appearing as the output of the Character Register are tested by $(Wb1 \oplus Wb2 \oplus Wb3)$. This parity logic term is true whenever there is an odd number of one's in a three bit group. It is used to reverse the state of Rwp for each such octal group containing an odd number of ones.

$$sRwp = \overline{W9} \overline{W4} \overline{Rwp} (Wb1 \oplus Wb2 \oplus Wb3) (T7 - T0) Qw1 + . . .$$

$$rRwp = \overline{W9} \overline{W4} Rwp (Wb1 \oplus Wb2 \oplus Wb3) (T7 - T0) Qw1 + . . .$$

The proper state of Rwp is thereby achieved, establishing whether the incoming character had a parity error. For example, if the following twelve bit character were received,

1 101 101 111 001


its parity bit would cause Rwp to start off by setting then being reversed to the reset state. Of the four octal groups, two of them have an odd number of ones. This would cause Rwp to change state twice as precession took place thus returning to the reset condition. With Rwp reset, no parity error would be reported by We.

3.62 The number of octal groups that must be checked for each character precession and which outputs of the Character Register to be monitored for 6-, 12-, and 24-bit characters must also be determined. The number of octal groups to check is solved by $(T7 - T0) Qw1$, which is true for 2, 4, or 8 pulse times in accordance with the character length. Table 3-2 lists the values of signal Qw1. The character length may be determined by hardwired logic or by gating of Wx12 and Wx24 signals (depending on the model). Which outputs to monitor is solved by using the same gating signals, Wx12 and Wx24, or by hardwiring to select the proper outputs. Whether gating or hardwiring is employed depends on the equipment model number. The variations in models are detailed in table 3-2.

3.63 The characters have now been read into the Character Register, precessed into the W-Register, and parity has been checked. The contents of the W-Register must now be transferred to the C-Register in the Central Processor Unit (CPU). Each time a WIM (AIM) instruction is executed, the contents of the W (A) and C-Registers are interchanged. At this time ($\emptyset 4$ of the WIM instruction execution) precession is blocked by $\overline{W4}$ and recirculation is blocked by Wxx. The condition $Wf \overline{W9}$ is used to signal the computer that the exchange can take place.

$$Wxx = Rx Pwy + . . . \quad (Pwy = 05 \text{ in Model } 92200)$$

Rx indicates that a WIM or MIW instruction is being processed. It is always false for the C and D channels, since no WIM or MIW type instructions exist for those channels.

$$Ww1 = \overline{W4} C21r (T7 - T0) Wxx$$

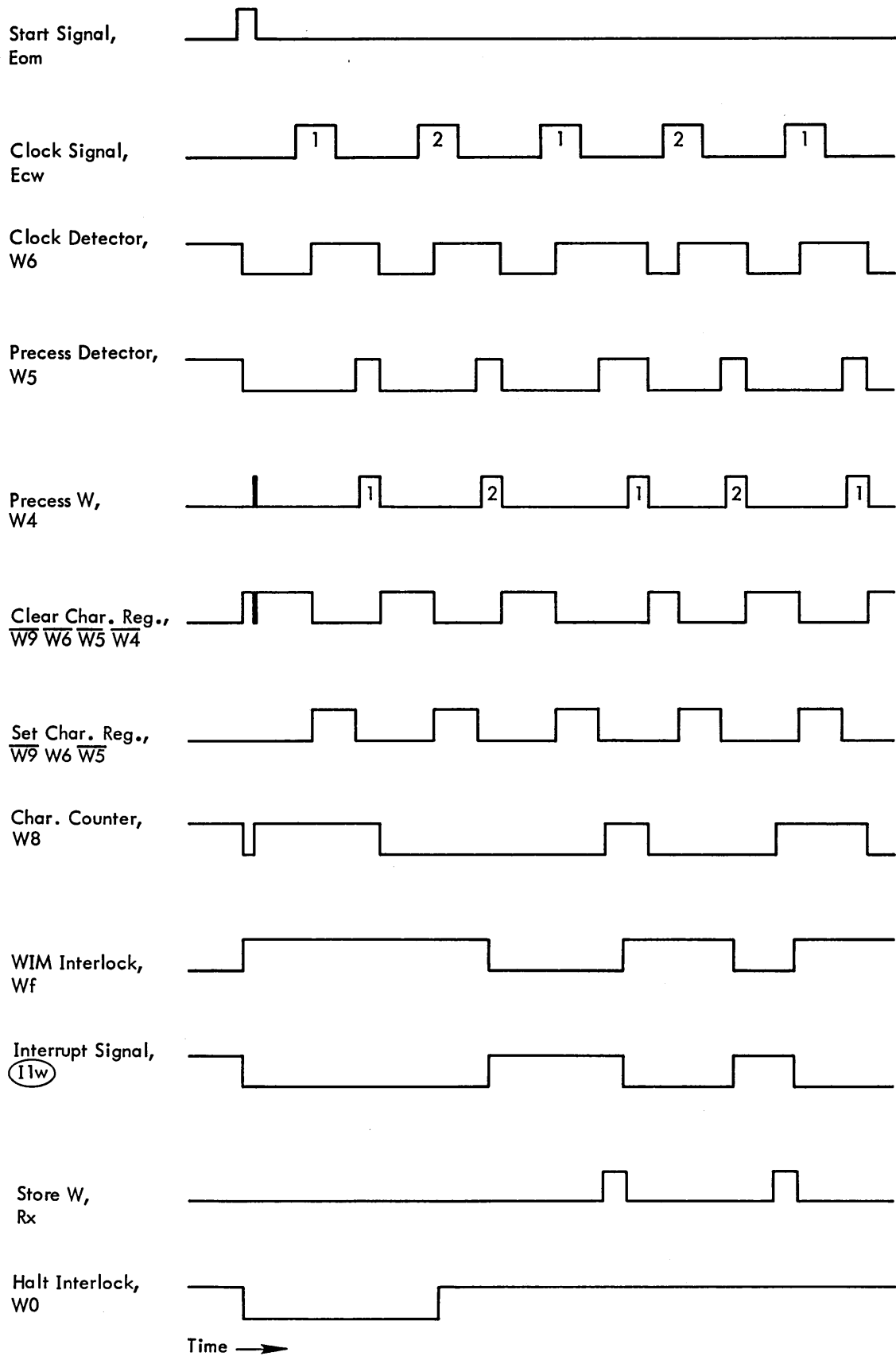


Figure 3-6. Input Timing Chart (Two Characters Per Word)

Table 3-2. Value of Parity Timing Signal Qw1

TMCC Model No.	Value of Qw1	Time When (T7 - T0) Qw1 is true
92200	$Qw1 = Qr1 \overline{Qr4} *$	T7, T6
92201	$Qw1 = Qr4 *$	T7, T6, T5, T4
92202	$Qw1 = 1$ (Qw1 is deleted) *	T7 through T0
93200/93221	$Qw1 = \overline{Wx12} \overline{Qr4} + \overline{Wx12} \overline{Wx24} \overline{Qr1} \overline{Qr4}$	T7, T6 if Wx12 and Wx24 are off. T7 through T4 if Wx12 is on. T7 through T0 if Wx24 is on.

*Qw1 is replaced by the signal shown.

$$Ww2 = \overline{W4} C22r (T7 - T0) Wxx$$

$$Ww3 = \overline{W4} C23r (T7 - T0) Wxx$$

The terms C21r, C22r, and C23r are logically equivalent to the outputs of the C-Register, C21, C22, and C23, respectively, but are implemented through special drivers for this transfer function in order to minimize the delay. The three parallel portions of the C-Register behave in a manner similar to one another, therefore only one is listed below (all three are illustrated in figure 3-7).

$$sC0 = Cr3 (Rx Rn1 + \dots)$$

$$rC0 = Cr3 \overline{(Rx Rn1 + \dots)}$$

$$sRn1 = Rwy1 \overline{Tsr} + \dots$$

$$rRn1 = \overline{Rwy1} \overline{Tsr} + \dots$$

3.64 The data transfer takes place during one machine cycle of the WIM instruction (Rx) and when completed, the computer Interlock Flip-Flop, Wf, is set to prepare the TMCC to accept the next input word.

$$sWf = Rx T0 Twy + \dots$$

Terminating an input operation can take place in one of several ways. The more sophisticated methods are

discussed in paragraph 3.131 dealing with the interlace. The simplest procedure is to program an EOM instruction to disconnect the peripheral unit after a sufficient number of WIM instructions have been processed. Other methods will now be discussed and illustrated as they apply to specific devices. Use will be made of the interlock and interrupt signals while considering these devices. These features are applicable to other I/O units as well.

3.65 A photoreader input process (refer to figure 3-8) can be terminated by detecting tape gap following a block of input data. The tape gap consists of one or more tape frames where only the sprocket hole is punched. However, the photoreader must be able to initially read through a tape gap or leader without terminating at every tape frame. This is accomplished by the Halt Interlock Flip-Flop, W0, which detects the start of a block of data and is used by the peripheral device coupler to inhibit sprocket clocks until the first (or second) character is sent to the TMCC. Clocks for these first characters are derived from the characters themselves. The equation for W0 is given below, but to fully understand its action, the equation for a typical photoreader clock signal is also shown.

$$sW0 = \overline{W9} W6 \overline{W8} Ecw + \dots$$

$$rW0 = Wc + \dots$$

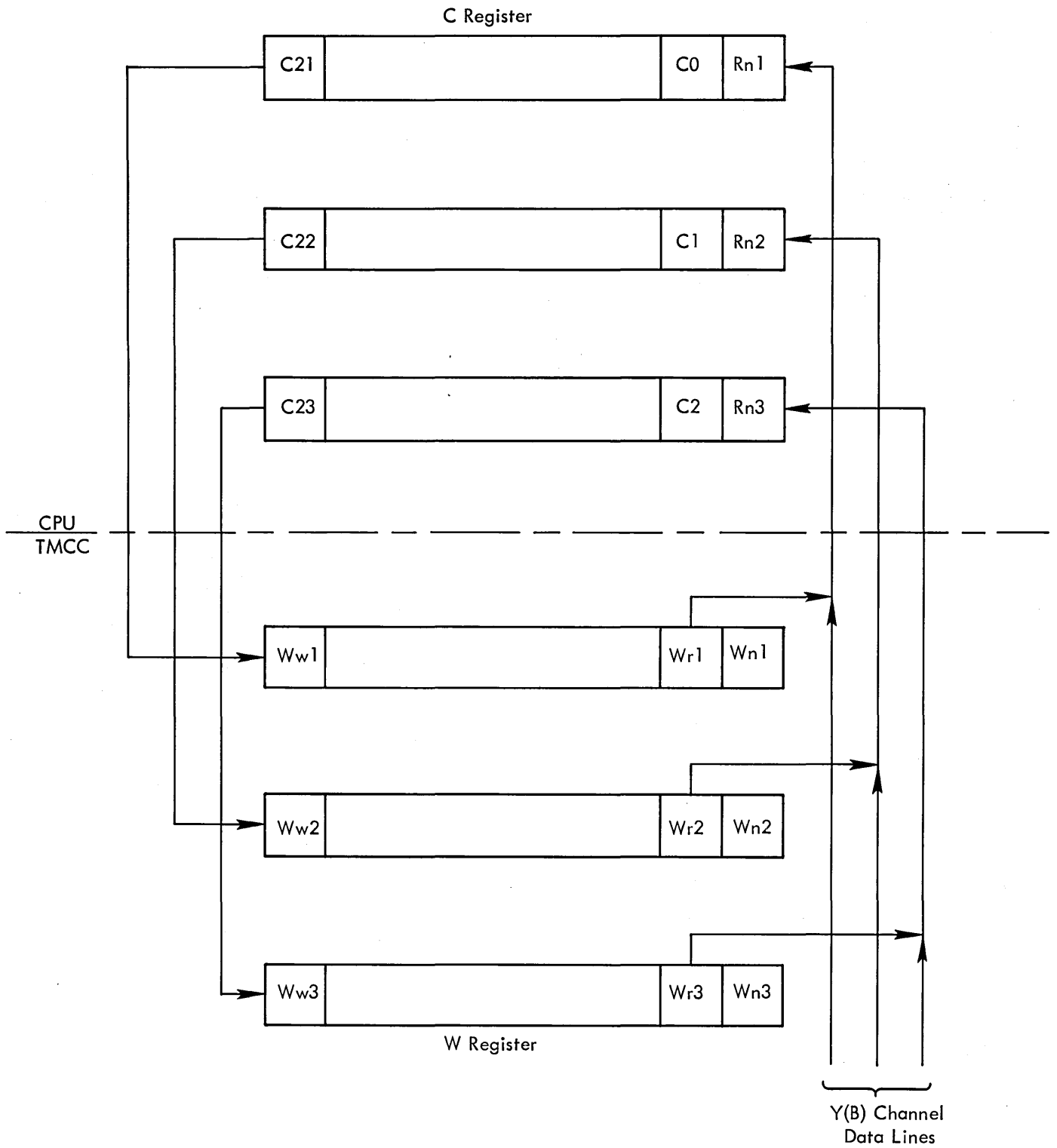


Figure 3-7. Data Transfer from W Register to C Register

$$E_{cw} = (Z_{w1} + Z_{w2} + Z_{w3} + Z_{w4} + Z_{w5} \\ + Z_{w6} + Z_{wp} + W_0) S_p R_e$$

In the latter equation, Z_{w1} through Z_{w6} are the character reader signals, Z_{wp} is the parity bit signal, S_p is the sprocket hole signal, and R_e is a photoreader enable level. The result of this combination is that while W_0 is reset, the clocks are derived from the character bits. After setting W_0 , the term in parentheses is always true so E_{cw} varies with S_p only. When the first character is read following leader, a clock is produced which sets W_6 . W_6 allows W_0 to be set. The term W_8 on the input of W_0 prevents setting of W_0 until the second character appears (in the four character per word mode). Waiting until the second character is desirable when reading from magnetic tape so a single noise character within a gap does not set W_0 . It is also important during a reverse scan of magnetic tape to avoid setting W_0 on the longitudinal parity character.

3.66 After W_0 has been set to enable the clocks to read in data, the End-of-Record Detector, W_g , is used to detect the first all-zeros, or gap, character (all zeros including the parity bit). W_g then enables the Halt Detector Flip-Flop, W_h .

$$sW_g = \overline{W_9} \overline{W_{10}} \overline{W_{11}} \overline{W_{12}} \overline{W_{13}} \\ (\overline{R_{w1}} \overline{R_{w2}} \overline{R_{w3}} \overline{R_{w4}} \overline{R_{w5}} \overline{R_{w6}} \overline{R_{wp}}) W_5 \\ (T_7 - T_0) (\overline{W_{10}} \overline{W_{11}} \overline{W_{12}} \overline{W_{13}} \overline{W_{14}})$$

$$rW_g = W_c$$

$$sW_h = W_g \overline{I_{wg}} T_8$$

The combination, $\overline{W_9} \overline{W_{10}} \overline{W_{11}} \overline{W_{12}} \overline{W_{13}}$, at the input of W_g decodes the fact that paper tape reader number 1 or number 2 is being used for input. The terms $(\overline{W_{10}} \overline{W_{11}} \overline{W_{12}} \overline{W_{13}} \overline{W_{14}})$ are redundant and are due to the logic mechanization.

3.67 After W_g sets, parity errors no longer set the Error Detector Flip-Flop, W_e .

$$sW_e = \overline{W_9} \overline{W_6} \overline{W_5} \overline{W_4} \overline{R_{wp}} \overline{W_g} \overline{N_{pw}} \overline{I_{wg}} + \dots$$

Disabling W_e avoids parity testing for the next few machine cycles. During this time a complete word is precessed into the W-Register if the input did not supply sufficient characters to finish the last word. In the case where W_h sets but the W-Register is not yet full, W_4 is set for successive cycles until the character counter reads 00. This process is termed "flushing". When W_7 W_8 read 00, W_f is reset and a W_c signal is generated to disconnect the TMCC.

$$sW_4 = W_h W_f T_8 + \dots$$

$$rW_4 = W_4 T_0 + \dots$$

$$sW_f = W_c \overline{W_h} + \dots$$

$$rW_f = \overline{W_7} \overline{W_8} W_4 (T_6 + T_5) + \dots$$

$$W_c = W_h \overline{W_f} (T_3 - T_0) + \dots$$

The clear signal, W_c , disconnects the tape reader and clears the character counter, End-of-Record Detector, and Halt Interlock.

$$rW_0 = W_c$$

$$rW_g = W_c$$

The Halt Detector is then reset at the next T_8 pulse time.

$$rW_h = W_h \overline{W_f} T_8$$

3.68 By allowing the character counter to count down to 00 before resetting, the final input word is filled-in with zero characters. The Phototape Termination Timing Charts illustrate the flow of this termination process. Figure 3-9 illustrates the case where the last input word does not contain a full complement of characters. Figure 3-10 illustrates the timing for the case where the last word does have all characters filled before an all zeros character is detected.

3.69 If the interrupt system is enabled during the termination process, the $\overline{W_f} W_h$ condition generates an interrupt signal to call for a final WIM instruction from the computer.

$$I_{2w} = (E_n + \textcircled{E_n}) \overline{I_{wg}} W_h \overline{W_f}$$

The signals E_n and $\textcircled{E_n}$ are programmable and manual enable signals for the interrupt system. The final WIM instruction, as provided by a halt subroutine, stores the last word even though it may not have originally contained a sufficient number of characters. The final WIM instruction may be executed before or after the channel is disconnected. However, additional WIM instructions after that cause the computer to lock-up because it would be waiting for another input word.

3.70 When operating with devices other than punched paper tape, the input process is very similar to that just described, although the method used to derive the clock signals may vary somewhat from one device to another. Also, the gap (End-of-Record) signal, when required, is normally supplied by the external device rather than developed within the TMCC as is done for paper tape.

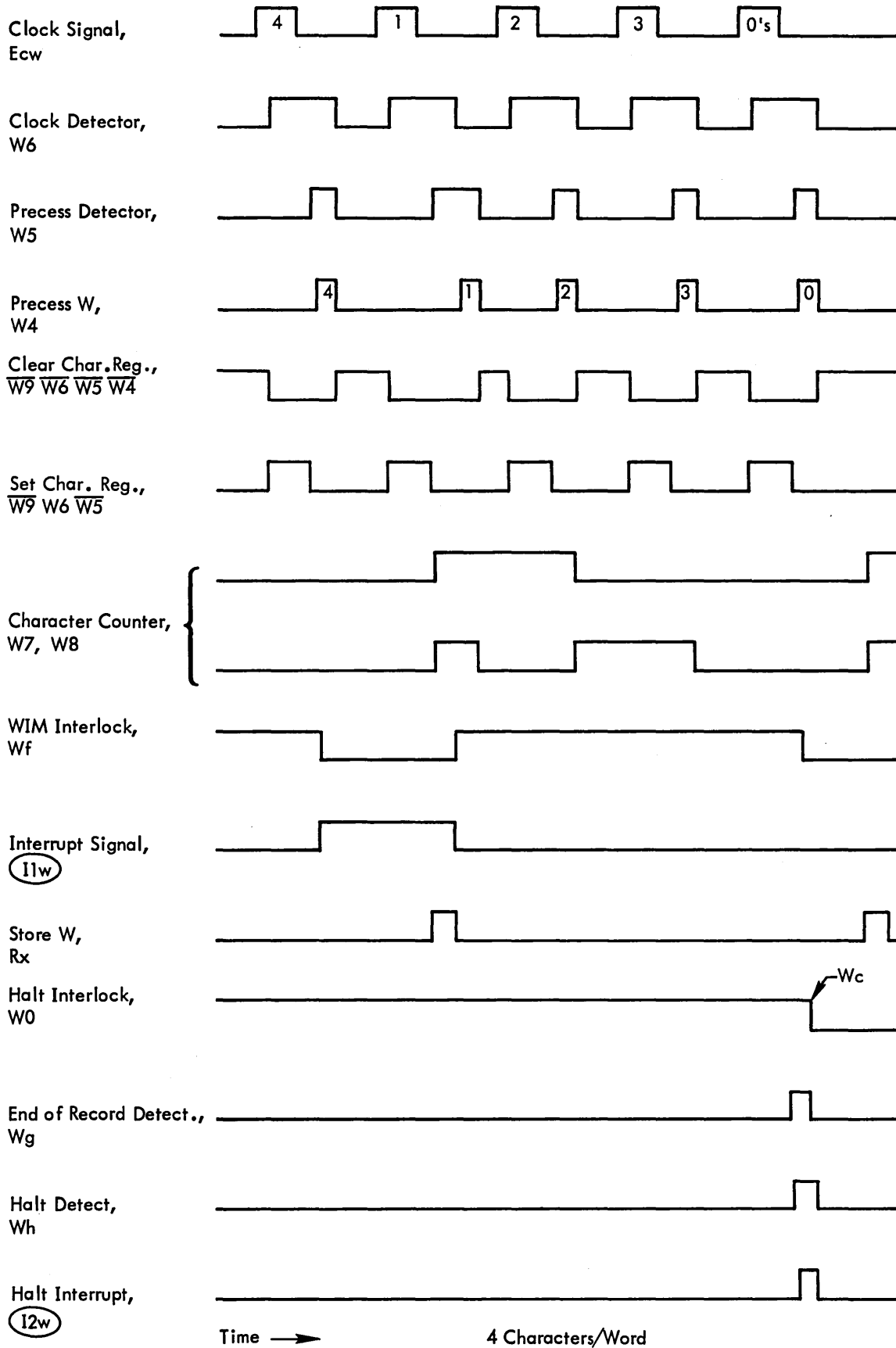


Figure 3-9. Termination Timing A. Phototape Input

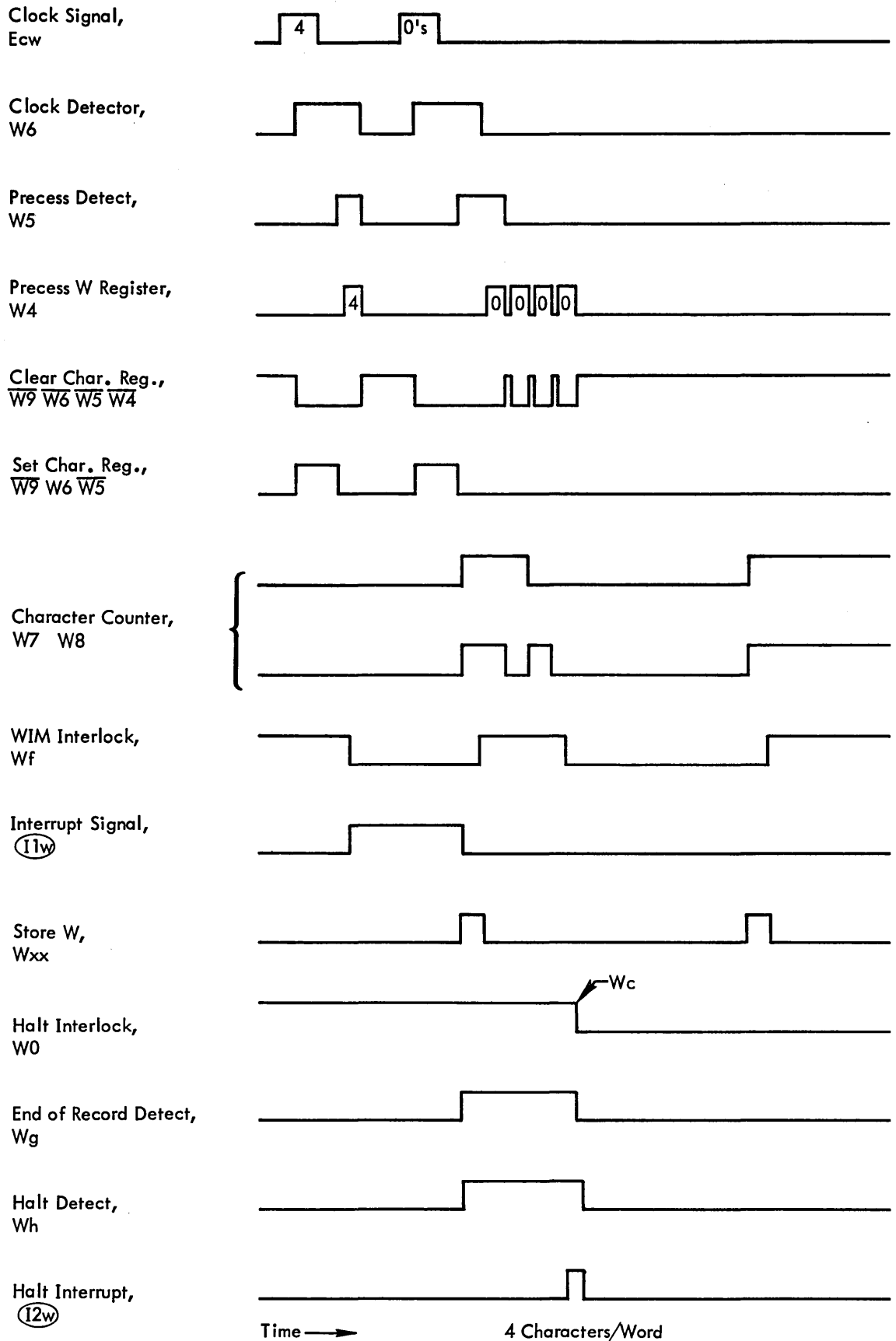


Figure 3-10. Termination Timing B - Phototape Input

3.71 An information flow diagram for magnetic tape input is illustrated in figure 3-11. The input timing for magnetic tape using the two character per word mode is illustrated in figure 3-12. When terminating, the magnetic tape unit generates a halt signal with a time delay triggered by the tape gap and W0. The halt detector is then triggered by the delayed signal.

$$sWg = Whs (\overline{T7 - T0}) \overline{W10 W11 W12 W13 W14}$$

$$rWg = Wc$$

$$sWh = Wg \overline{Iwg} T8$$

Whs is a halt signal from the peripheral device. In this case, it is the delayed signal from the magnetic tape unit. The delay allows time for the tape reader to check the longitudinal parity character following the block of data. Refer to figure 3-11. When a longitudinal parity error occurs, the tape reader sends an error signal \overline{Wes} and the Error Detector is set. The \overline{Wes} signal may also result from a rate error.

$$sWe = Wes + . . .$$

Other devices may also supply error inputs via the \overline{Wes} line. Unless inhibited by Npw, character parity is checked by Rwp when precession takes place.

3.72 Input termination timing for magnetic tape is illustrated in figure 3-13. The figure illustrates the case in which the input furnishes only three characters for the last four character word and the remaining character is filled with zeros. Also illustrated in figure 3-13 are the Halt Interrupt, calling for one more WIM instruction, and the final WIM instruction, itself.

3.73 Another type of input operation in which the TMCC participates is the scanning of magnetic tape. The process is similar to the usual magnetic tape reading process except that the character counter is prevented from reaching 00 again after W0 is set. The counter flip-flop, W8, is held in the "one" or set state by:

$$sW8 = \overline{W7 W9 W10 W11 Wh} + . . .$$

In this equation, $\overline{W9 W10 W11}$ indicates a scan process has been programmed. A forward scan can be initiated by an EOM 0363X instruction. It may also be programmed by modifying a read process. While magnetic tape is being read, if an EOM14000 instruction is given, the flip-flop, W10, is set to convert directly from reading to scanning of the same tape.

$$sW10 = (Ioc C12 \overline{C17} \overline{C19} \overline{C20} \overline{C21} \overline{C22} \overline{C23}) \overline{W9 W10} + . . .$$

With W8 of the character counter held in a set state, the Buffer Full Signal, $\overline{W9 Wf}$, and the normal interrupt signals are prevented by keeping Wf from being reset.

$$rWf = \overline{W7 W8} W4 (T6 + T5) + . . .$$

$$I1w = \overline{Wf W0 Wh} (En + \overline{En}) \overline{Iw} \overline{Ew} \overline{Iwg} + . . .$$

$$I2w = (En + \overline{En}) \overline{Iwg} Wh \overline{Wf} + . . .$$

This allows each input character to be precessed into the W-Register without requiring WIM instructions to set Wf. However, when the end of a data block is reached, Wf is reset by a gap signal from the tape reader.

$$rWf = \overline{W9 W10 W11 W0} Mtgw \overline{W7} (T6 + T5) Wh$$

And an interrupt is generated.

$$I1w = \overline{Wf W0 Wh} (. . .)$$

This interrupt calls on the computer to enter a subroutine to execute a WIM instruction. The WIM instruction stores the last four characters read from the tape and which were precessed into the W-Register. The WIM instruction sets W7 and Wf in preparation for scanning another record.

$$sW7 = Wxx Wn1 (\overline{T7 - T0}) W4 + . . .$$

$$sWf = Tx T0 Pwy + . . . \quad (Pwy = 05 \text{ for } 92200)$$

3.74 Based on a block counting program or an examination of the last four characters, the computer may reset W0 with another EOM0363X instruction, to cause the scan process to continue without a pause through the gap and into the next record. If W0 is not reset, the scan process is terminated by a \overline{Whs} signal from the tape unit in a manner similar to that previously described for normal input termination. When the scan process does terminate, a halt interrupt signal is generated.

$$I2w = (En + \overline{En}) \overline{Iwg} Wh \overline{Wf}$$

3.75 During the scan process the character parity is checked by the Rwp flip-flop and the longitudinal parity is checked by the tape unit (only if W0 is reset after the longitudinal parity character). Any error during the scan process sets the Error Detector flip-flop, We.

3.76 A reverse scan of magnetic tape is started by an EOM 0563X instruction. The process is similar to a forward scan, except that the WIM instruction at the end of a data block scan stores the first four characters of the block in reverse order. The longitudinal parity is not properly checked, with the result that We may be erroneously set.

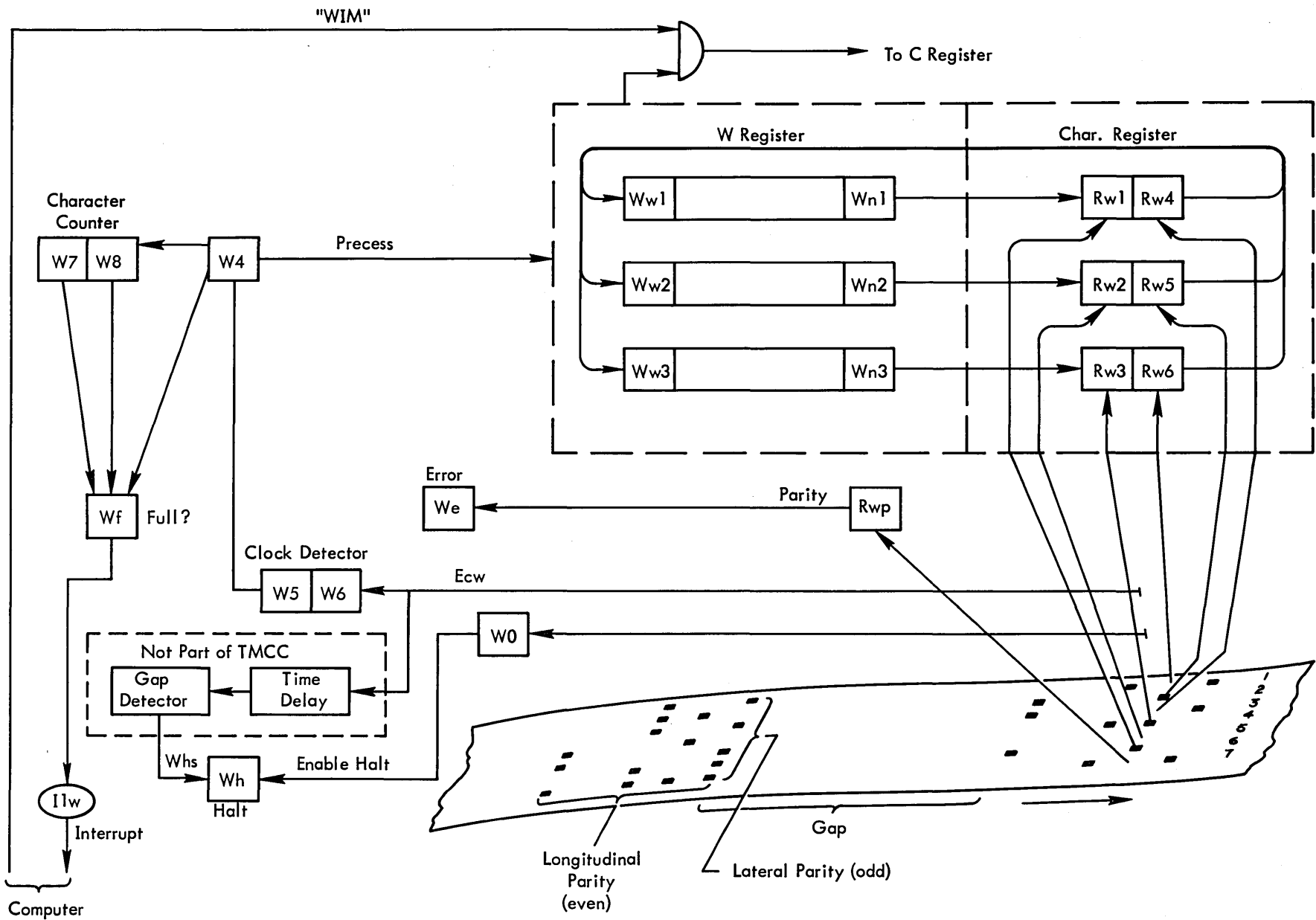


Figure 3-11. Information Flow Diagram - Magnetic Tape

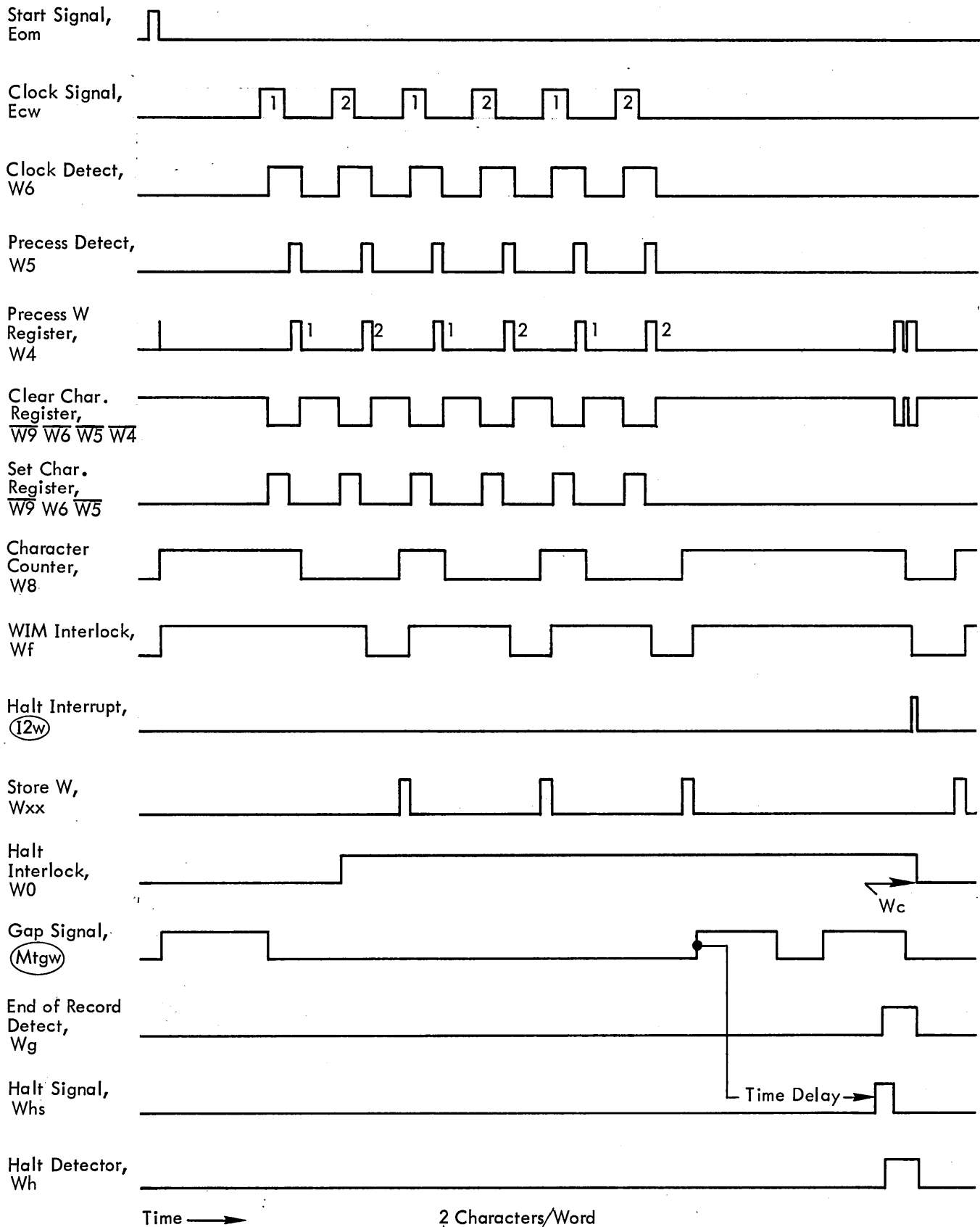


Figure 3-12. Input Timing - Magnetic Tape

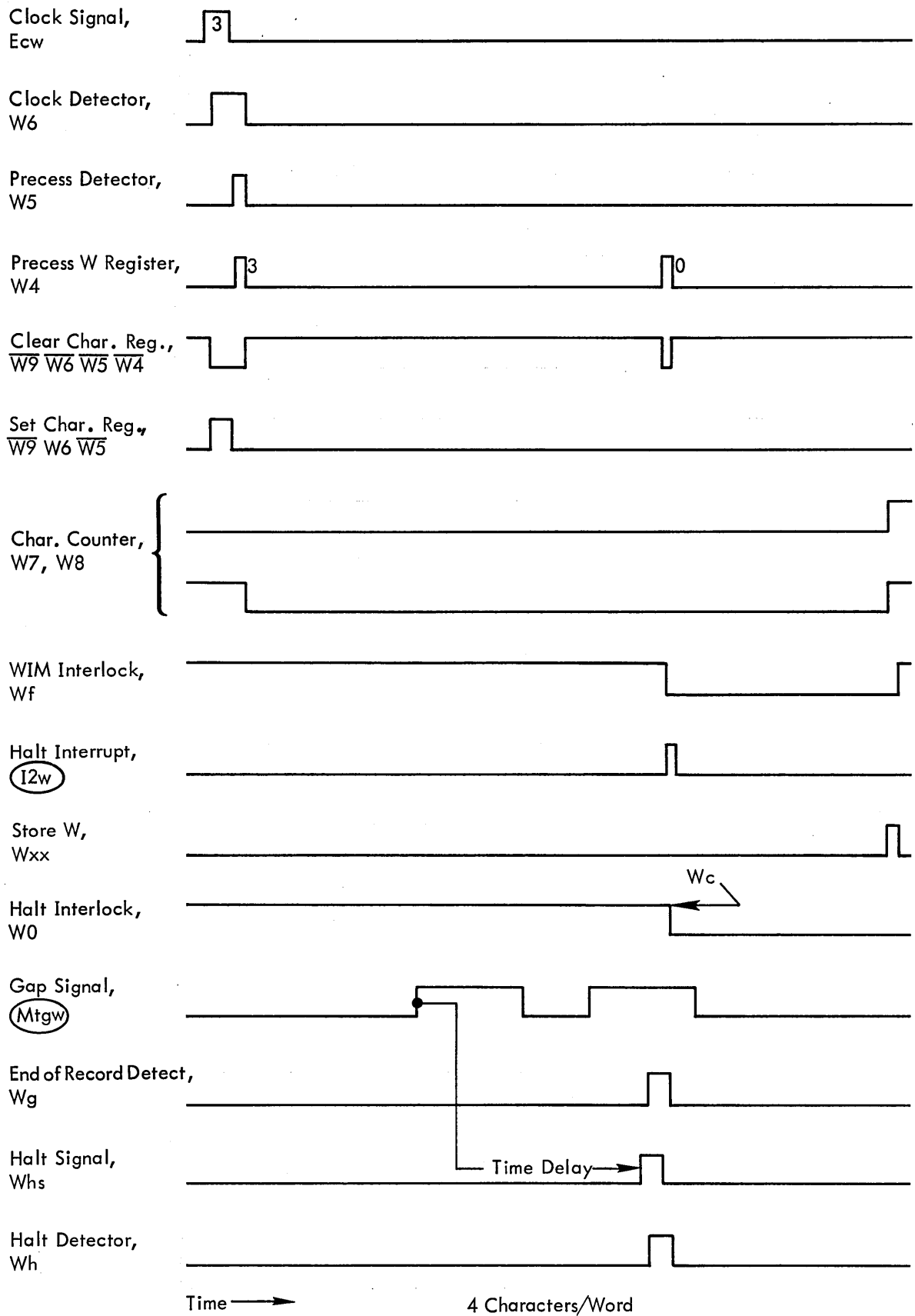


Figure 3-13. Input Termination Timing - Magnetic Tape

3.77 In the Magnetic Tape Forward Scan Timing Chart (refer to figure 3-14) and Reverse Scan Timing Chart (refer to figure 3-15), when the first interrupt informs the computer that the gap has been reached, a WIM instruction stores the last four characters read. Another EOM instruction then resets W0 to continue the scan process through the next block

$$rW0 = (\text{loc } C12 \overline{C17} \overline{C19} \overline{C20} \overline{C21} \overline{C23}) \\ W9 T0 + \dots$$

Dotted lines indicate the reaction if the process were terminated. In case of termination, W0 allows the Tape Reader Halt Signal, Whs, to come through.

$$sWg = Whs \overline{T7 - T0} \overline{W10} \overline{W11} \overline{W12} \overline{W13} \overline{W14} + \dots \\ sWh = Wg \overline{Wg} T8 + \dots$$

Whs sets Wg, which in turn sets Wh. After Wh is set, the Character Counter Flip-Flop, W8, is no longer held on and a count-down begins.

$$sW8 = \overline{W7} \overline{W9} W10 W11 \overline{Wh} + \dots \\ rW8 = W8 W4 T0 + Wc$$

This process is very similar to that described for terminating a normal read operation and Wf is reset as soon as the count-down reaches a point where W7 and W8 are both reset.

$$rWf = \overline{W7} \overline{W8} W4 (T6 + T5)$$

A clear signal, Wc, then resets the TMCC after the four precessions.

$$Wc = Wh \overline{Wf} (T3 - T0) + \dots$$

3.78 OUTPUT PROCESS (W9 true)

3.79 When an EOM0XXXX or EOM4XXXX instruction is executed to start an output process, an interlock signal, \overline{Wf} ($W0 + \dots$) is immediately sent to the computer. And, if enabled, an interrupt calls for the computer to load the W-Register with the first output word. If the interrupt is disabled, the program should supply the MIW loading instruction before it is needed (i. e. before a clock signal is received from the peripheral unit).

$$rWf = Ws C18 + \dots \text{ (denotes W register is empty)} \\ sW0 = Ws C18 W9 + \dots$$

$$\textcircled{11w} = \overline{Wf} W0 \overline{Wh} (En + \textcircled{En}) \overline{Iw} \overline{Ew} \overline{Iwg} + \dots$$

3.80 Each MIW instruction, or time-share operation if interlace is being used (W_{xx}), exchanges the W-Register and the C-Register. Refer to figure 3-16.

$$Ww1 = \overline{W4} C21r (T7 - T0) W_{xx} + \dots$$

$$Ww2 = \overline{W4} C22r (T7 - T0) W_{xx} + \dots$$

$$Ww3 = \overline{W4} C23r W_{xx} + \dots$$

$$sC0 = Cr3 \left[(\overline{Tsr} + Rx) Rn1 + \dots \right]$$

$$rC0 = Cr3 \left[(\overline{Tsr} + Rx) Rn1 + \dots \right]$$

similarly for C1 and C2

$$sRn1 = Rwy1 \overline{Tsr} + Rr1 Tsr$$

$$rRn1 = \overline{Rwy1} \overline{Tsr} + \overline{Rr1} Tsr$$

similarly for Rn2 and Rn3

$$W_{xx} = Rx Pwy + \dots \quad (Pwy = 05 \text{ for } 92200, \\ Rx \text{ is always false for } \\ \text{TMCC-C})$$

3.81 As with the input process, $\overline{W_{xx}}$ blocks recirculation and with W4 false precession is blocked at this time also. When precession does take place, as enabled by W4, an output character is shifted from the W-Register to the Character Register and Rwp is used to generate an odd parity bit. During the output operation the parity flip-flop operates in a manner similar to its action during input.

3.82 Rwp is initially set by Wf W5 T8. Then, when precession occurs, each octal group coming from the W-Register is examined for an odd or even number of ones. This checking is done by $(Wn1 \oplus Wn2 \oplus Wn3)$ which is true whenever there are one or three ones in Wn1, Wn2, and Wn3. Each time this term is true during the checking period, Rwp is switched to its opposite state.

$$sRwp = W9 W4 \overline{Rwp} \overline{W_{xx}} (Wn1 \oplus Wn2 \\ \oplus Wn3) Qw2 (T7 - T0) \\ + Wf W5 T8 \overline{Rwp} + \dots$$

$$rRwp = W9 W4 Rwp \overline{W_{xx}} (Wn1 \oplus Wn2 \\ \oplus Wn3) Qw2 (T7 - T0) + \dots$$

Qw2 is a term similar to the Qw1 used during the input process. Qw2 with $(T7 - T0)$ establishes the pulse times during which the parity checking is done. It may either be hardwired or function with the gating signals Wx12 and Wx24 depending on the equipment model number. Refer to table 3-3.

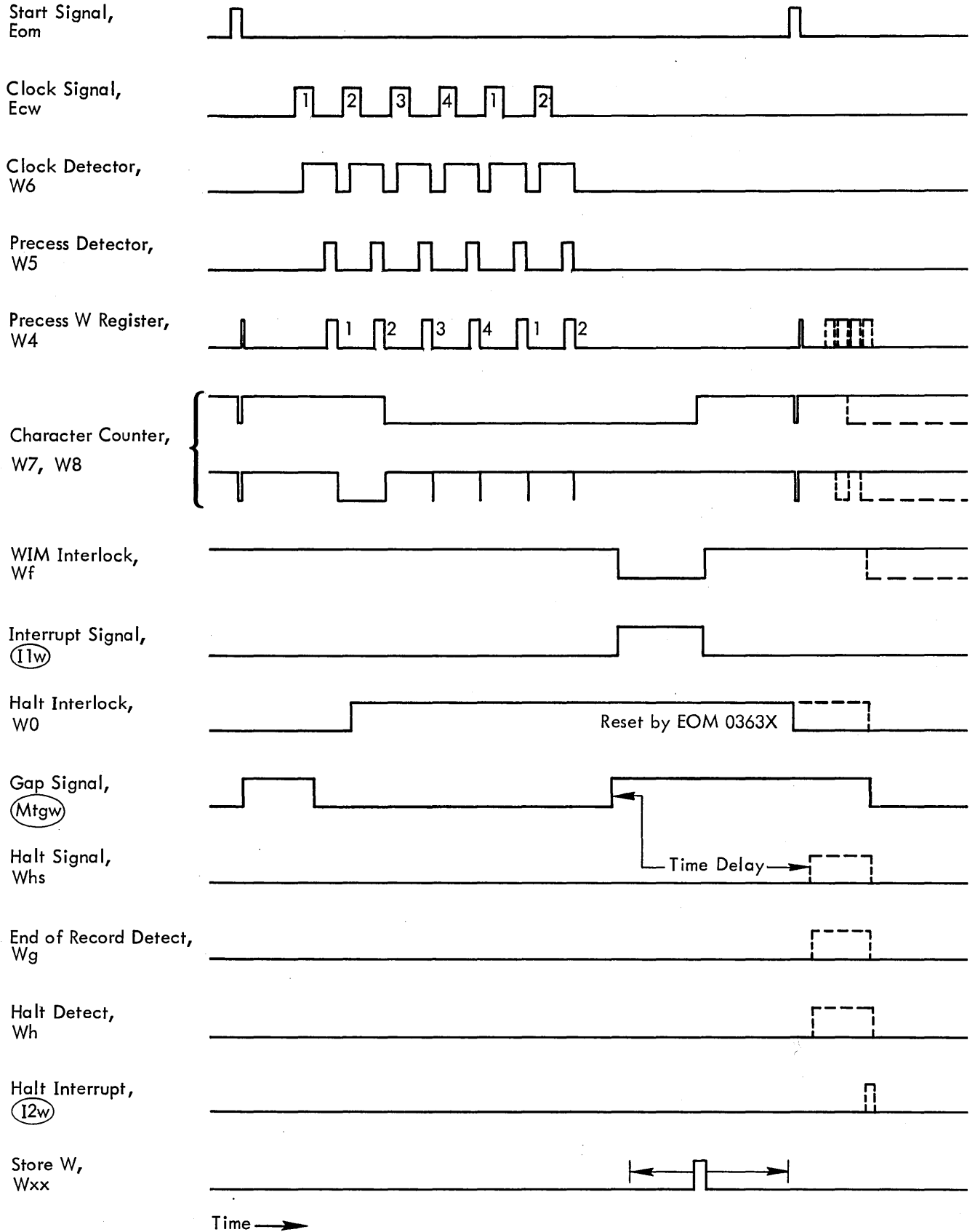


Figure 3-14. Forward Scan Timing Chart - Magnetic Tape

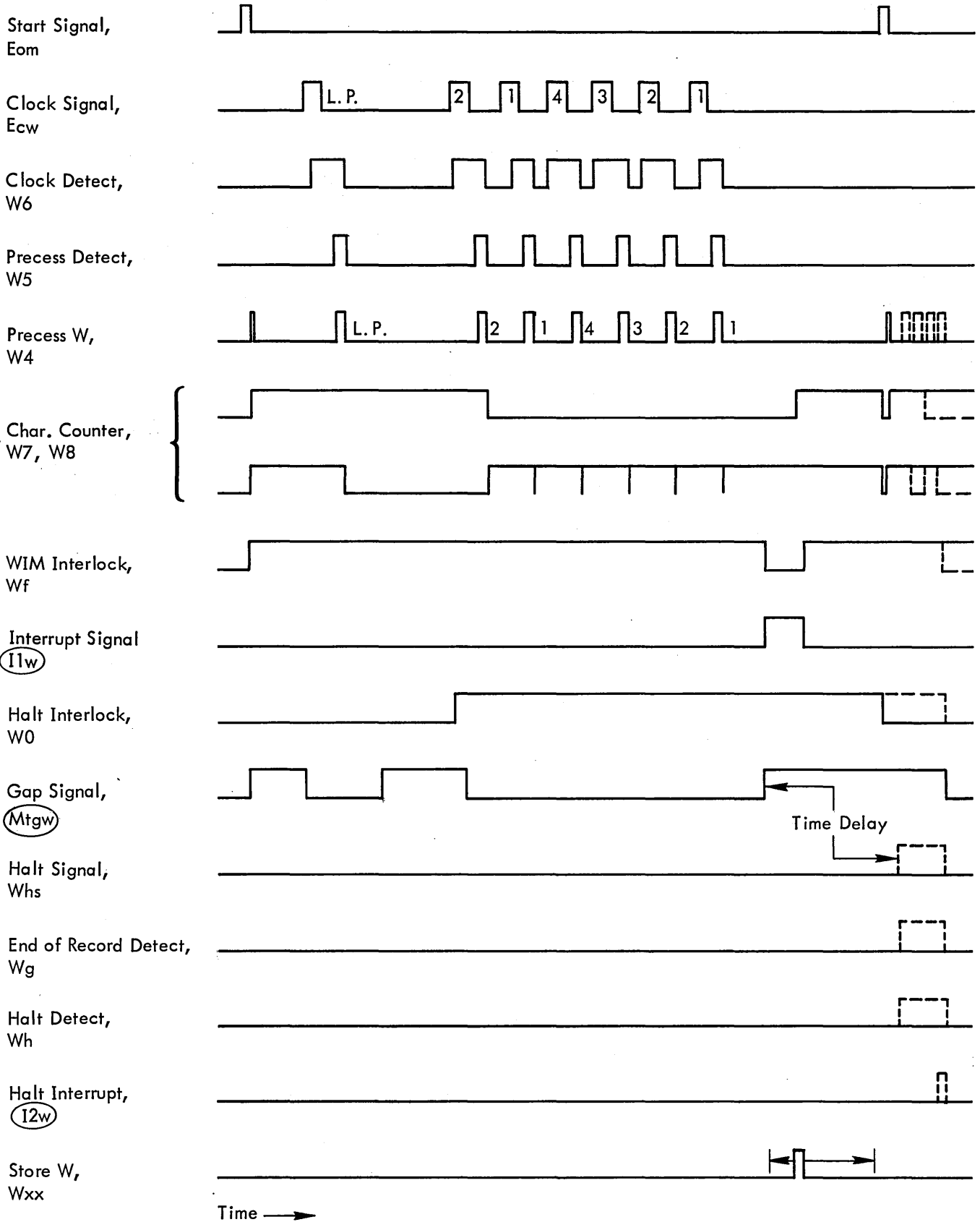


Figure 3-15. Reverse Scan Timing Chart - Magnetic Tape

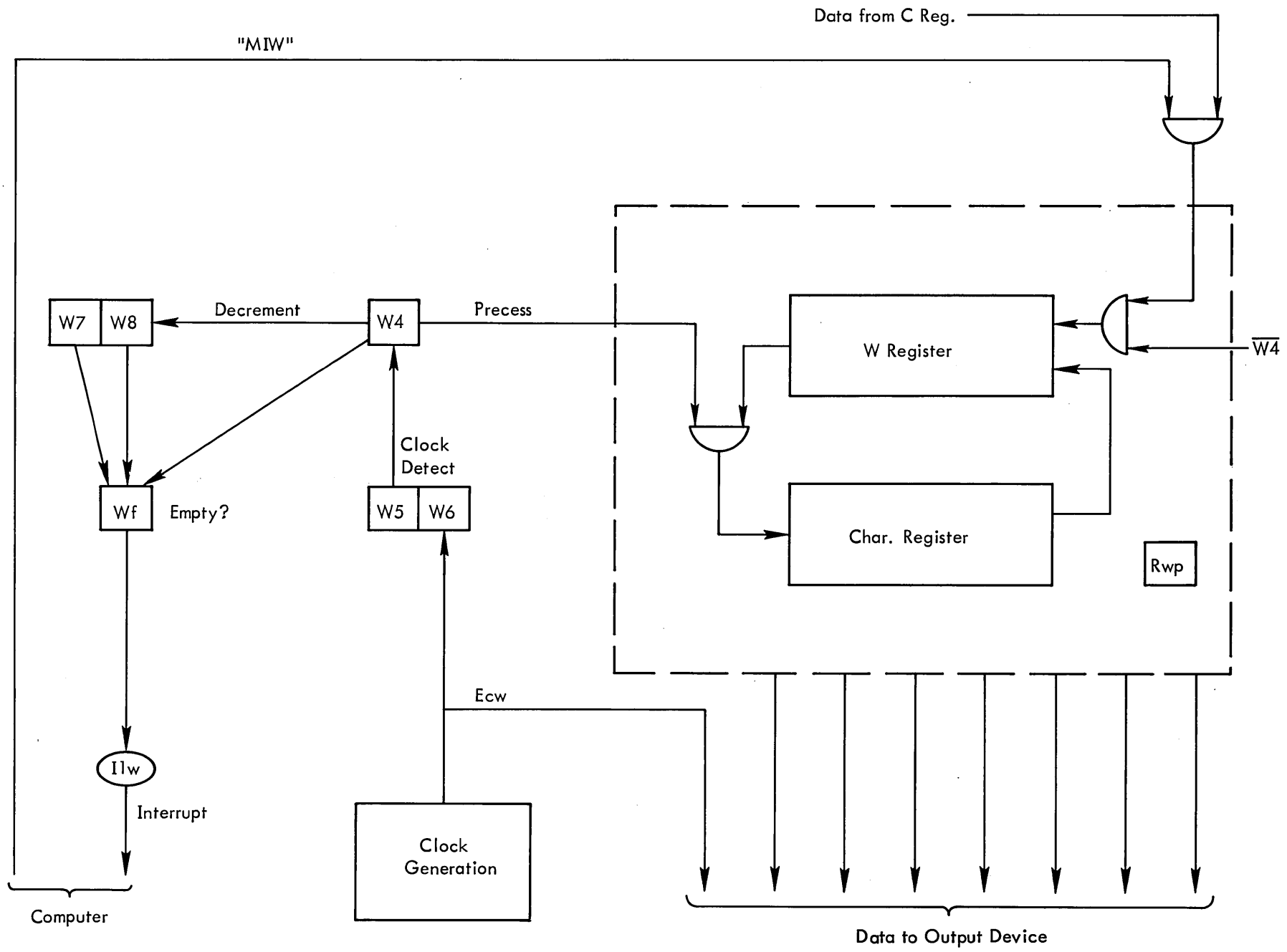


Figure 3-16. Output Information Flow Diagram

Table 3-3. Value of Qw2 Parity Timing Signal

TMCC Model No	Value of Qw2	Pulse Times When (T7 - T0) Qw2 is True
92200	$Qw2 = Qr2 \overline{Qr4} *$	T0, T1
92201	$Qw2 = \overline{Qr4} *$	T3 through T0
92202	$Qw2 = 1$ (Qw2 is deleted)	T7 through T0
93200/93221	$Qw2 = \overline{Wx12 Qr4 + Wx12 Wx24 Qr2 Qr4}$	T0 and T1, if Wx12 and Wx24 are both off. T3 through T0, if Wx12 is on T7 through T0, if Wx24 is on

*Qw2 is replaced by the signal shown above for Models 92200/201/202.

3.83 As an example of output parity generation, consider the twelve-bit character 111 010 101 001. Rwp would start in the set condition and would then switch states three times, once for each of the octals containing an odd number of ones. Thus, Rwp would conclude its switching operation in the reset state to produce a zero parity bit for an odd parity output character.

3.84 The basic flow of the output process is illustrated in figure 3-17. The execution of the second MIW instruction is shown occurring late to depict how the Character Register is cleared when a new output character is not available. If another output character is still not available when the next clock signal, Ecw appears, the error detector (We) is set as is done during an input process.

$$sWe = W0 \overline{W6} W5 Ecw T8 + \dots$$

$$rWe = Wc \overline{Wh}$$

3.85 In the output processes (as in the input processes), the Character Register is initially cleared by $\overline{W9} \overline{W6} \overline{W5} \overline{W4}$. The first output character to be read is, therefore, all zeros. This is appropriate for a leader or gap in paper tape punching. However, for some forms of output, such as typing or leaderless punching, and magnetic tape writing (the tape unit automatically generates leader), the first output character should be in the Character Register before the first clock signal. For this type of output, an EOM2XXXX instruction with a "one" bit in C13 is used. This code bit is used to set W5 which then causes the first loading of the W-Register to be followed directly by precession of the first output character into the Character Register. In this case the sequence of operations again starts by

resetting Wf and setting W0 with the resulting interrupt Ilw. However, now W5 is also set by the Ws signal. Then, as usual, when the MIW instruction is executed and Rx goes true, Wf is set. With W5 on, W4 can be set at pulse time T8 just after the data transfer takes place. Figure 3-18, Output Timing Chart 2, illustrates the flow of this process.

$$sWf = Rx T0 Pwy + \dots \quad (Pwy = 05 \text{ for } 92200)$$

$$rWf = Ws C18 + \dots$$

$$sW5 = Ws C13 C18 \overline{W10} \overline{W11} \overline{W12} \overline{W13} \overline{W14} + \dots$$

$$rW5 = W4 T0 + Wc$$

$$sW4 = Wr Wf T8 \overline{Wg} + \dots$$

$$rW4 = W4 T0 + \dots$$

3.86 To terminate an output process, the MIW instruction which loads the last output word into the W-Register is followed by an EOM14000 instruction to reset W0. An EOM14100 instruction is used to terminate the Y channel.

$$rW0 = (Ioc C12 \overline{C17} \overline{C19} \overline{C20} \overline{C21} \overline{C22} \overline{C23})$$

$$W9 T0 + \dots$$

$$Ioc = Ioc1 \overline{C1} \overline{Er} Qr3$$

$$Ioc1 = Eom \overline{C10} C11$$

\overline{Er} is a signal inhibiting Eom and Ioc when the Interlace Prepare flip-flop (to be discussed later) has been set. The C1 bit, which appears on Ioc, indicates a W or Y channel instruction. C1 is used with the C and D channels.

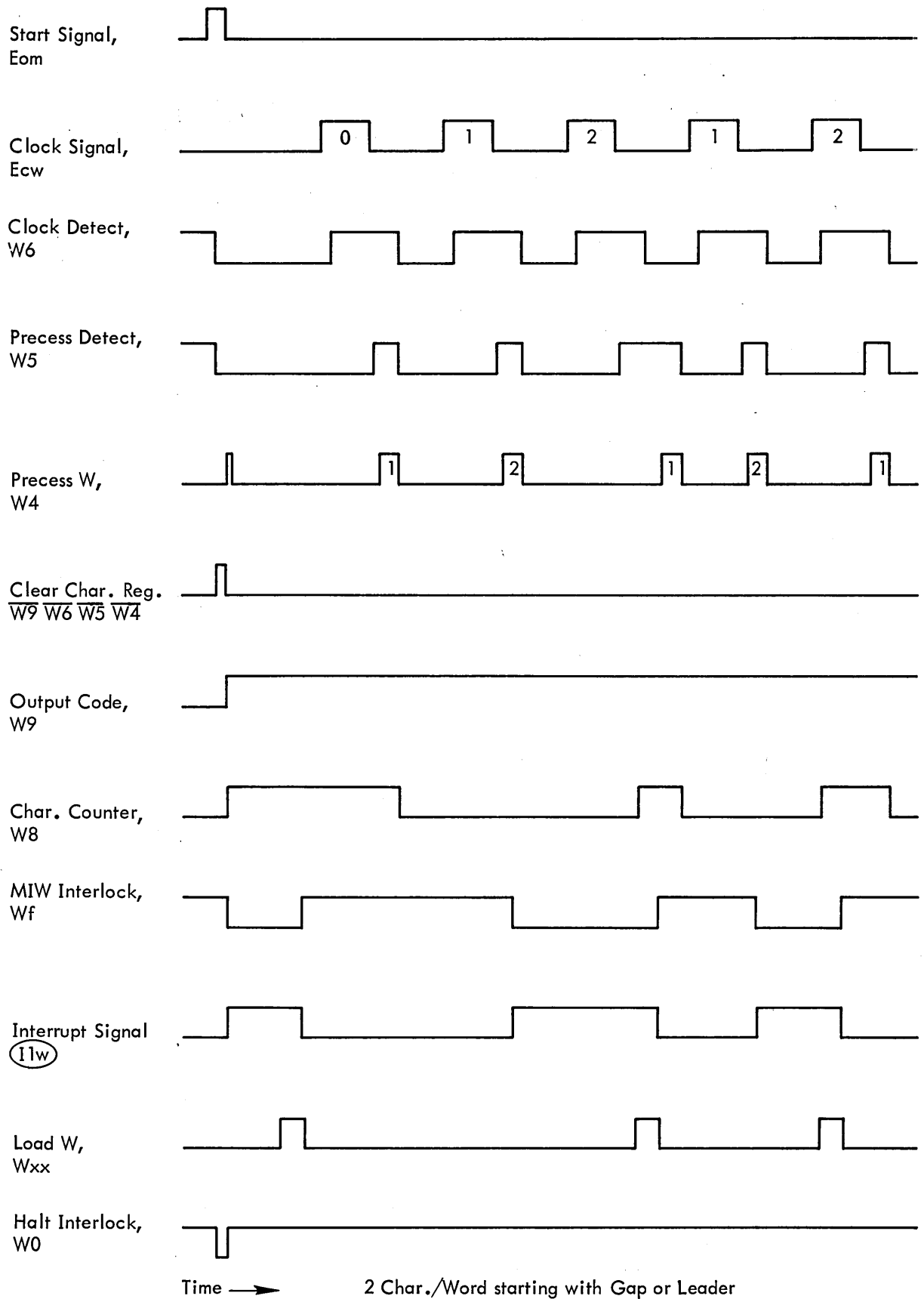


Figure 3-17. Output Timing Chart 1

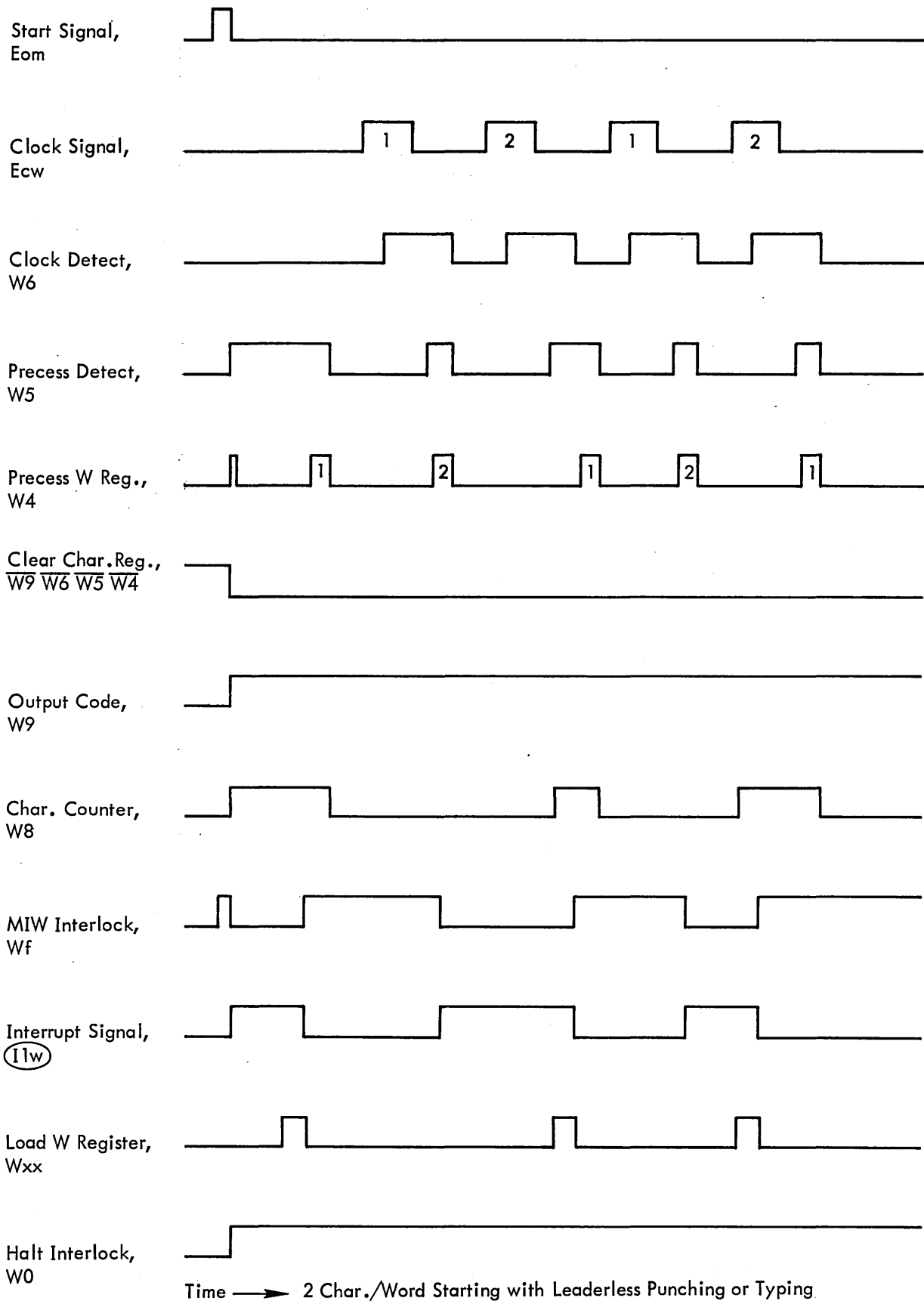


Figure 3-18. Output Timing Chart 2

3.87 Resetting of W0 results in the following: Further normal interrupt signals are blocked.

$$I_{w1} = \overline{Wf} \overline{W0} \overline{Wh} (. . .)$$

Further late-load error signals are blocked.

$$s_{We} = W0 \overline{W6} W5 E_{cw} T8 + . . .$$

The WIM/MIW Interlock Signal Wf W0 is blocked. As a result of inhibiting the interlock and interrupt signals, no further MIW instructions are processed and, therefore, Wf is not set again after the last character is precessed into the Character Register.

$$s_{Wf} = R_x T0 P_{wy} + . . . \quad (P_{wy} = 05 \text{ for } 92200)$$

$$r_{Wf} = \overline{W7} \overline{W8} W4 (T6 + T5) + . . .$$

With Wf reset, W4 is prevented from setting after the last output character is precessed. This results in the state, $\overline{W0} \overline{W4} W5 \overline{W6}$, following the precession.

$$s_{W4} = W5 Wf T8 \overline{Wg} + . . .$$

$$r_{W4} = W4 T0 + W4 T8$$

$$s_{W6} = \overline{W5} E_{cw} T8 \overline{W10} \overline{W11} \overline{W12} \overline{W13} \overline{W14} + . . .$$

$$r_{W6} = W5 W0 + Wc$$

$$s_{W5} = \overline{W5} W6 \overline{E_{cw}} T0 + . . .$$

$$s_{W5} = W4 T0 + Wc$$

The state, $\overline{W0} W5 \overline{W6}$, is used to set the halt detector unless magnetic tape is being used ($\overline{W11}$ indicates not magnetic tape).

$$s_{Wh} = W9 \overline{W11} \overline{W0} W5 \overline{W6} (\overline{I_{wg}} + . . .) T8 + . . .$$

When the output is to a magnetic tape unit, the state $\overline{W0} W5 \overline{W6}$ is sent to that unit and after a suitable delay (while the tape unit generates a longitudinal parity character) a Whs halt signal is received back to set Wh.

$$s_{Wh} = Whs W11 T8 + . . .$$

Regardless of the method of setting Wh to terminate an output process, the Halt Interrupt Signal, $\overline{I_{2w}}$, is generated in the cycle in which Wh is set.

$$I_{2w} = (E_n + \overline{E_n}) \overline{I_{wg}} Wh \overline{Wf} + . . .$$

And a clear signal is generated.

$$Wc = Wh \overline{Wf} (T3 - T0) + . . .$$

Then Wh is reset, as well as the rest of the TMCC.

$$r_{Wh} = Wh \overline{Wf} T8 + Wc (T6 + T5)$$

The Output Termination Timing Charts indicate the flow of these output termination processes for devices other than magnetic tape (refer to figure 3-19) and for magnetic tape (refer to figure 3-20).

3.88 Two additional timing charts are included to illustrate the output sequences for specific devices. The flow of the complete output process for a paper tape punch operation is illustrated in figure 3-21. The EOM0XX4X start instruction causes tape leader to be punched while the device inhibits clock signals. An all zeros character is also punched for the first output clock signal. After the last output character is processed, a halt interrupt signal is generated.

3.89 The flow of an output process using magnetic tape is illustrated in figure 3-22. A time delay triggered by the EOM02X5X start instruction causes a tape gap to be recorded first while inhibiting output clock signals. $\overline{W0} W5 \overline{W6}$ signals the tape unit to count three clocks and record the longitudinal parity character, and triggers a second time delay to cause a gap to be recorded after the data block. When the gap is completed, the tape unit generates a \overline{Whs} signal to halt the output process. Each character parity and the longitudinal parity of the characters reproduced at the read head are checked by the tape unit and an error signal, \overline{Wes} , is generated to set We for any detected errors. Several other error conditions are also checked, such as slew, amplitude, and rate error.

$$s_{We} = Wes + . . .$$

3.90 A special case of a magnetic tape output operation is the erase function. The tape erase is started by an EOM01X7X instruction. The erase procedure is performed in the same way as any other output to magnetic tape but the W10 bit is used by the tape unit to cause writing of all zero data regardless of what may be appearing at the character register outputs.

3.91 SYS GATE

3.92 The system control EOM instruction has little effect on the TMCC. It is included here only because the Sys signal is gated through the TMCC channel.

3.93 When an EOM instruction, containing "ones" in bits 10 and 11, is executed, an Sys signal is generated as an output from the TMCC on the Sys line.

$$Sys = E_{om} C10 C11 \overline{C9}$$

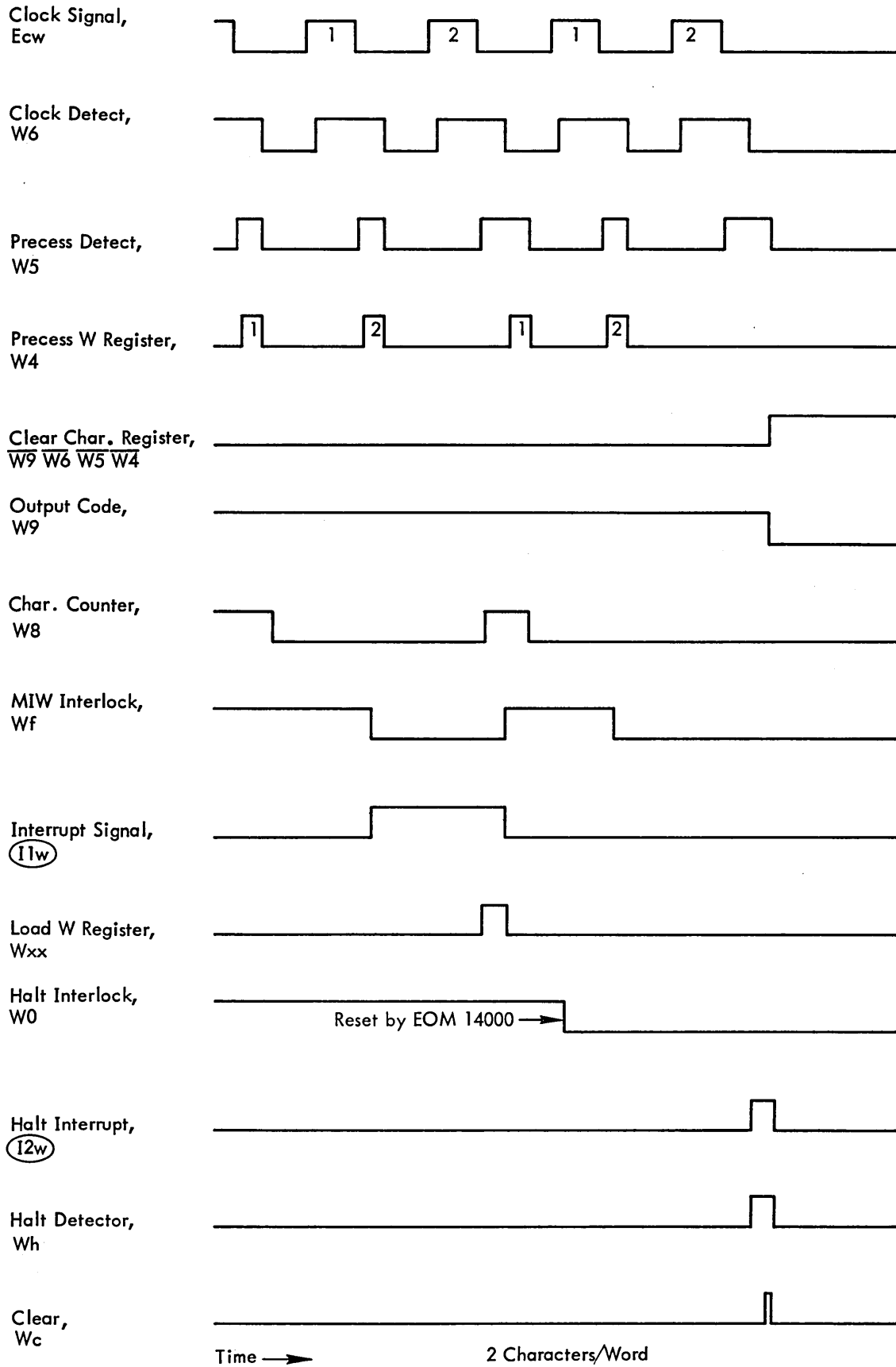


Figure 3-19. Output Termination Timing (Except Magnetic Tape)

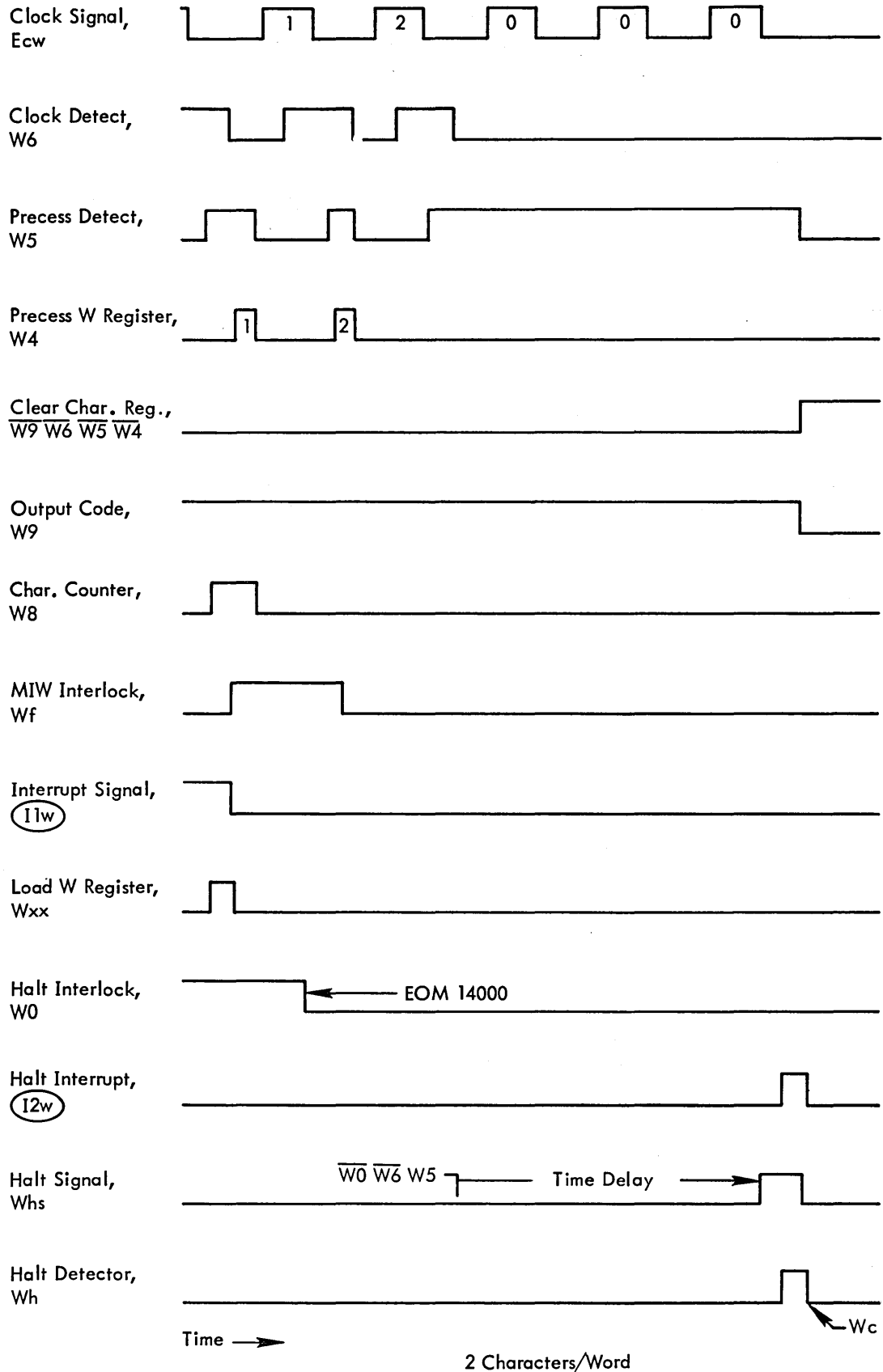


Figure 3-20. Output Termination Timing - Magnetic Tape

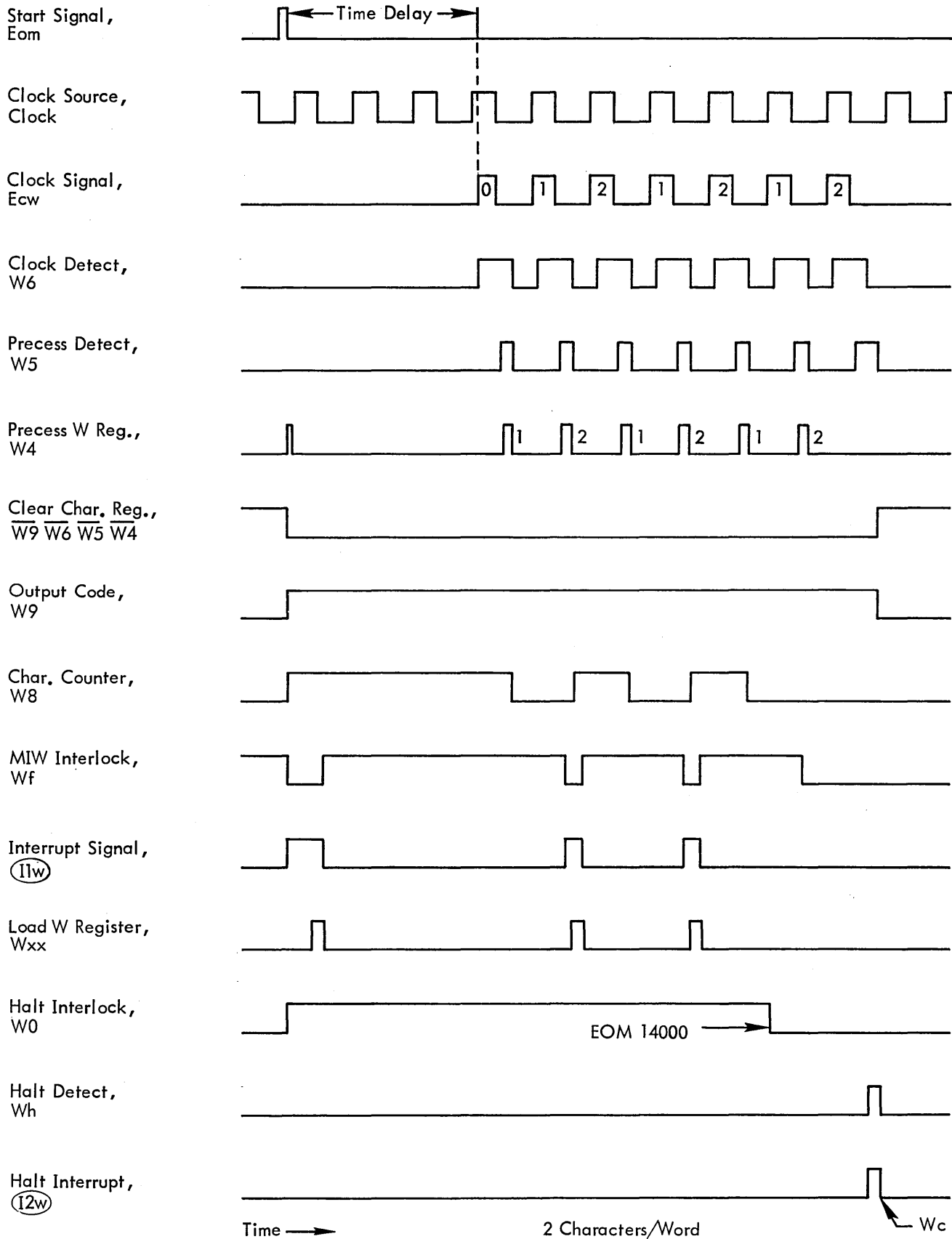


Figure 3-21. Output Timing Chart - Punch

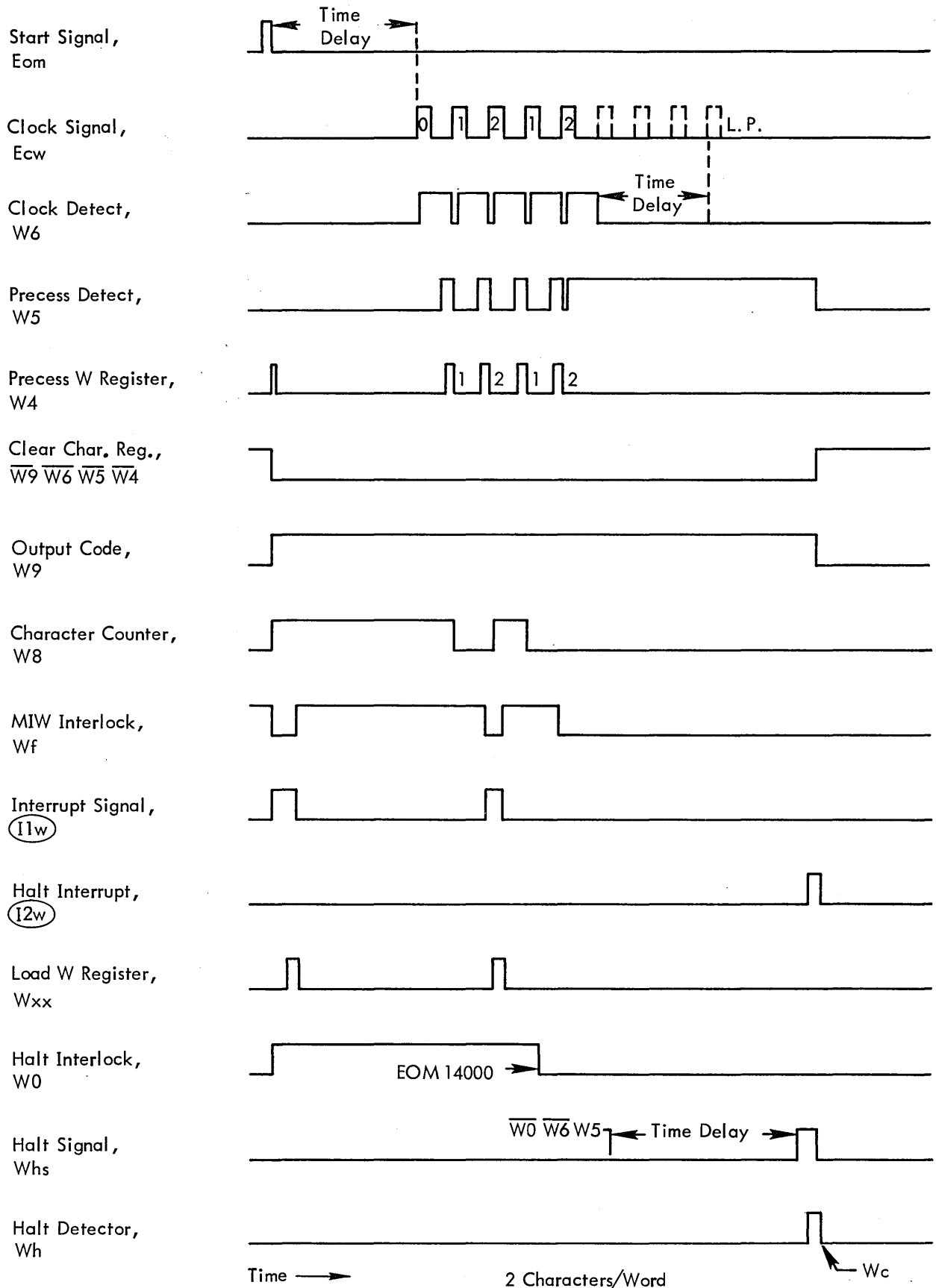


Figure 3-22. Output Timing Chart - Magnetic Tape

$$\text{Sys} = \overline{\text{Sys}} (\text{T5} - \text{T1}) \overline{\text{Tsr}} \text{Pwy}$$

(delete Pwy for 92200 and 92210)

The Sys signal may be used in conjunction with the C-Register output lines C0 through C24 to select and control special system devices. One Sys line is provided on each TMCC W (A) channel or each W(A) + Y(B) channel pair. One more Sys line is provided on the C or C + D channels (for a maximum of two lines).

3. 94 SKIP GATE

3. 95 The Skip Gate, Skr , and its associated line driver, $\overline{\text{Skrz}}$, provide interrogation response signals for use with SKS instructions.

The equation for the line driver is:

$$\overline{\text{Skrz}} = \overline{\text{Skr}} (\overline{\text{C10}} \overline{\text{C11}} \overline{\text{Ssc}}) (\overline{\text{C1}} \overline{\text{C9}} \overline{\text{C10}} \overline{\text{C11}} \text{Sio}) \dots *$$

* $\overline{\text{C1}}$ becomes C1 for TMCC-C and TMCC-D

This signal is inverted in the CPU and sampled during execution of an SKS instruction to determine if a response has been received. At times other than when being tested, the signal is changing levels and has no meaning. The inverse of $\overline{\text{Skrz}}$ is:

$$\text{Skrz} = \text{Skr} + \text{C10} \text{C11} \text{Ssc} + \overline{\text{C1}} \overline{\text{C9}} \overline{\text{C10}} \overline{\text{C11}} \text{Sio} + \dots$$

3. 96 The dashed lines indicate that other response signals may be connected to satisfy special system requirements. Sio is a response from the addressed peripheral I/O unit. The addressing is performed using SKS instruction bits 18 through 23 and the same address codes as assigned for EOM instructions. The output lines C0 through C24 provide the necessary addressing and control connections to the peripheral units. As seen above, Sio is used in conjunction with $\overline{\text{C1}} \overline{\text{C9}} \overline{\text{C10}} \overline{\text{C11}}$. Bits 1 and 9 (along with bit 17 decoded by the peripheral device) select one channel out of four TMCC and four DACC channels. Bits 10 and 11 determine the type of the SKS instruction.

3. 97 Ssc is an interrogation response from external system equipment. This type of interrogation is selected by an SKS instruction with "ones" in bits 10 and 11. The control and address bits may be assigned as required. Again the $\text{C0} - \text{C24}$ lines provide the necessary connections.

3. 98 The remaining term in Skrz is Skr , the output of the Skip-Gate in the TMCC. The use of Skr is similar

to that of Sio just described except the SKS instruction address bits (19-23) are "zeros". This causes the instruction to address the I/O channel itself rather than peripheral devices.

$$\begin{aligned} \text{Skr} = & \overline{\text{C1}} \overline{\text{C17}} \overline{\text{C9}} \overline{\text{C10}} \overline{\text{C11}} \overline{\text{C19}} \overline{\text{C20}} \overline{\text{C21}} \overline{\text{C22}} \overline{\text{C23}} \text{C15} \text{Wsc} \\ & + \overline{\text{C1}} \overline{\text{C17}} \overline{\text{C9}} \overline{\text{C10}} \overline{\text{C11}} \overline{\text{C19}} \overline{\text{C20}} \overline{\text{C21}} \overline{\text{C22}} \overline{\text{C23}} \text{C12} \\ & \cdot \overline{\text{W10}} \overline{\text{W11}} \overline{\text{W12}} \overline{\text{W13}} \overline{\text{W14}} \\ & + \overline{\text{C1}} \overline{\text{C17}} \overline{\text{C9}} \overline{\text{C10}} \overline{\text{C11}} \overline{\text{C19}} \overline{\text{C20}} \overline{\text{C21}} \overline{\text{C22}} \overline{\text{C23}} \text{C13} \text{Iwf} \\ & + \overline{\text{C1}} \overline{\text{C17}} \overline{\text{C9}} \overline{\text{C10}} \overline{\text{C11}} \overline{\text{C19}} \overline{\text{C20}} \overline{\text{C21}} \overline{\text{C22}} \overline{\text{C23}} \text{C14} \overline{\text{We}} \end{aligned}$$

3. 99 In these equations, bits 1, 9, and 17 select one of eight channels (TMCCs and DACCs). Bits 10 and 11 determine the type of SKS instruction. "Zeros" in bits 19 through 23 address the channel rather than a peripheral unit. Bits 12, 13, 14, and 15 select the particular test function. A "one" in bit 15 tests the Signal-Complete flip-flop. Similarly, a "one" in bits 12, 13, or 14 tests the Unit Address Register, the Interlace Address Register, or the Error Flip-Flop, respectively.

3. 100 Two additional TMCC channel tests, similar to two of those just described, are provided by the Skr gate. These also test the Unit Address Register and the Error Flip-Flop but with a different type of SKS instruction. This is done to provide program compatibility with the SDS 910/920 Computers.

$$\begin{aligned} \text{Skr} = & \text{C10} \overline{\text{C11}} \text{C14} \overline{\text{W10}} \overline{\text{W11}} \overline{\text{W12}} \overline{\text{W13}} \overline{\text{W14}} \overline{\text{C1}} \\ & + \text{C10} \overline{\text{C11}} \text{C20} \overline{\text{We}} \overline{\text{C1}} + \dots \end{aligned}$$

3. 101 The equations shown in this discussion are specifically applicable to the W(A) channel TMCC. To select the other channels C1 must replace $\overline{\text{C1}}$ for channels C and D. To distinguish between W and Y or between C and D, C17 is switched (except in the last two equations where C14 and C20 are used for W and are changed to C13 and C19 for the Y channel).

3. 102 INTERLACE, COMPATIBLE MODE (Time Share)

3. 103 The Interlace register is enabled by a "one" in bit position 9 of either BUC or IOC type of EOM instruction. These produce Buc or Ioc signals which are derived from EOM0XXXX and EOM1XXXX instructions, respectively. Therefore, the interlace register can be enabled by the same instruction that sets up the other TMCC registers, (i. e., by Buc), or it can be enabled without disturbing the rest of the TMCC, (i. e., by Ioc). Either of the EOM instructions clears the entire interlace register then sets the Enable Flip-Flop, Ew .

Interlace Clear: $\text{Iwc} = \text{Eom} \text{C9} \overline{\text{C10}} \overline{\text{C1}} \overline{\text{C17}} (\text{T3} - \text{T0})$
 $\overline{\text{C1}}$ Becomes C1 for TMCC-C

Interlace Prepare: $sEw = Iwc \overline{Ew} (T3 - T0)$
 $rEw = Wc T0 + \dots$

3.104 For a Buc instruction, Ew is first reset by Wc at pulse time $T0$ then reset during the next pulse time. The above equation for Iwc is specifically related to the $W(A)$ channel. To select any of the other three TMCC channels, different combinations of $C1$ and $C17$ are used.

Channel	C1	C17
W(A)	0	0
Y(B)	0	1
C	1	0
D	1	1

With Ew set, the computer can preset the interlace register with the starting memory address and the word count for an I/O process. Refer to figure 3-23. The loading of these counts is accomplished by a POT instruction. The POT instruction, which produces a $Pot1$ signal, is ordinarily used to parallel transfer data from the C-Register to an external device. $Pot1$, then, produces the loading signal,

$$Iwp = Pot1 (T6 + T5) Ew$$

then resets Ew ,

$$rEw = Pot1 (T3 - T0) Ew + \dots$$

3.105 The Word Counter, $Wc0$ through $Wc14$, is initially cleared by Iwc which sets all stages of the counter. The one's complement of the count is then produced by resetting the counter flip-flops for corresponding "ones" in the C-Register during Iwp . Actually, this applies to only the ten least significant bits of the counter. If a word count greater than 1023 is needed, the five most significant bits of the counter must be preset by a second EOM instruction. This instruction must be executed after the Interlace is enabled and before the POT instruction is executed. An Ioc instruction, without bit 9, is given for this purpose and it generates an Iwe load signal.

$$Iwe = Ioc1 (T6 + T5) Ew$$

Iwe then resets the most significant bits of the Word Counter to complete the storage of the one's complement count.

$$sWc14 = Iwc + \dots$$

$$rWc14 = Iwp C9 + \dots$$

$$sWc5 = Iwc + \dots$$

$$rWc5 = Iwp C0 + \dots$$

$$sWc4 = Iwc + \dots$$

$$rWc4 = Iwe C23 + \dots$$

$$sWc0 = Iwc + \dots$$

$$rWc0 = Iwe C19 + \dots$$

3.106 Using the one's complement, the Word Counter can count up rather than down and produce a termination signal when it contains "ones" in all of its stages. The counting is performed by triggering each stage of the counter on the falling edge of a previous stage. For example:

$$sWc13 = \overline{Ew} \overline{Wc13} \underline{Wc14} \overline{Ew}$$

$$rWc13 = \overline{Ew} \underline{Wc13} \underline{Wc14} \overline{Ew}$$

where the underlined term represents the clock or triggering level. This method of counting conserves gating but produces a propagation delay through the counter. This delay is minimized to a satisfactory level with some increase in the number of gates by arranging the fifteen counter stages into five octal groups. The first stage of each group is connected to only the last stage of the previous group. An example of one octal group is:

$$sWc10 = \overline{Ec} \overline{Wc10} \underline{Wc11} \overline{Ew}$$

$$rWc10 = \overline{Ec} \underline{Wc10} \underline{Wc11} \overline{Ew}$$

$$sWc9 = \overline{Ew} \overline{Wc9} \underline{Wc10} \underline{Wc11} \overline{Ew}$$

$$rWc9 = \overline{Ew} \underline{Wc9} \underline{Wc10} \underline{Wc11} \overline{Ew}$$

$$sWc8 = \overline{Ew} \overline{Wc8} \underline{Wc9} \underline{Wc10} \underline{Wc11} \overline{Ew}$$

$$rWc8 = \overline{Ew} \underline{Wc8} \underline{Wc9} \underline{Wc10} \underline{Wc11} \overline{Ew}$$

Here it is seen that for this octal group, a common clock term $\underline{Wc11} \overline{Ew}$ is used for all three stages.

3.107 This reduces the propagation delay for each group to the delay that would be expected for a single flip-flop. The delay for the entire fifteen stage counter is therefore equivalent to that of five flip-flop stages. The double appearance of \overline{Ew} is due to the clocking arrangement and the flip-flop module layout. During the loading operations, Ew is set to allow the clock pulses to appear as the $Gc4$ clock for use by the flip-flop. After Ew is reset and the counting is to take place, $Gc4$ is held off while $\underline{Ew} \underline{Wc14}$ performs the triggering. Refer to figure 3-24.

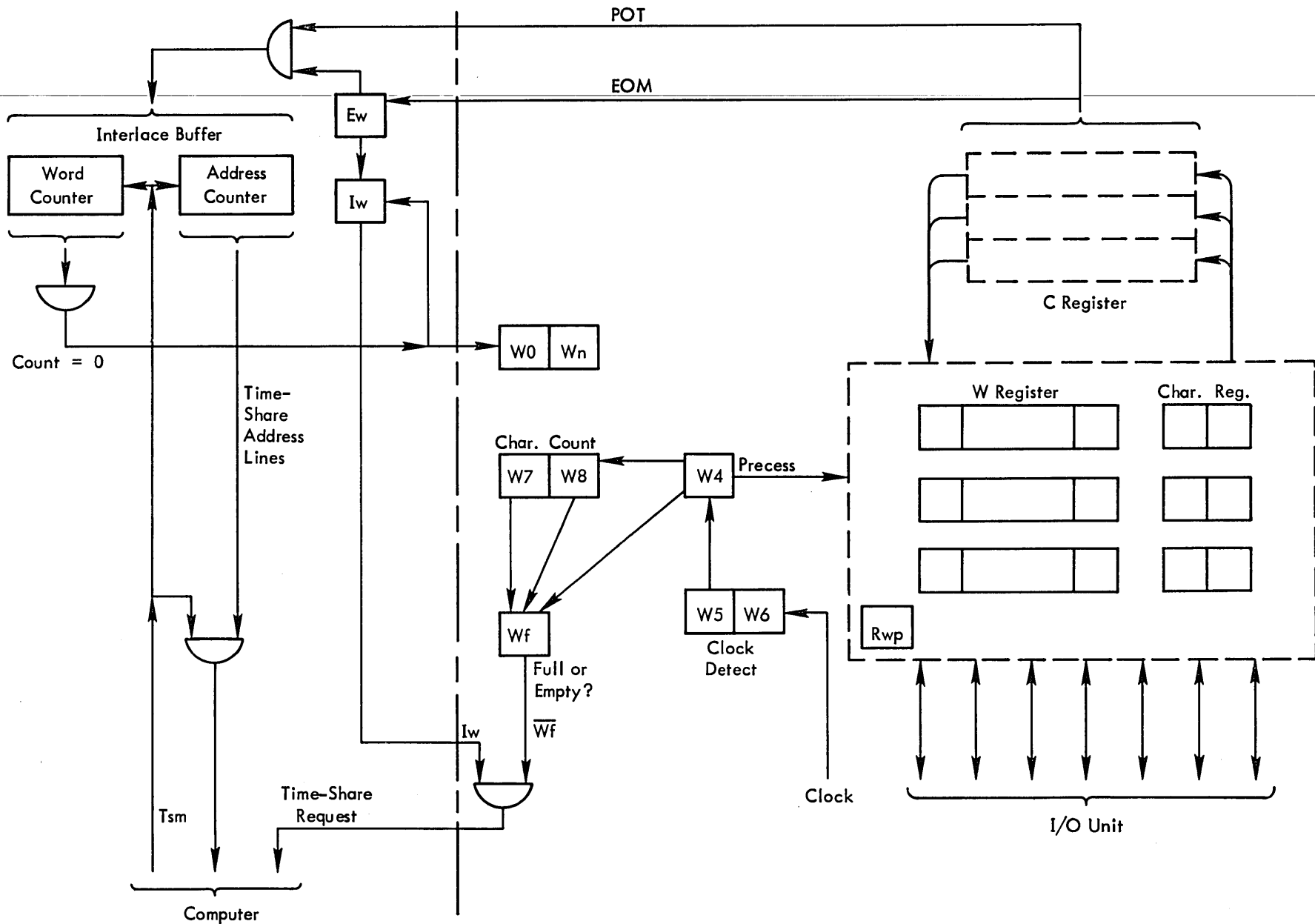


Figure 3-23. Information Flow - Interlace Operation

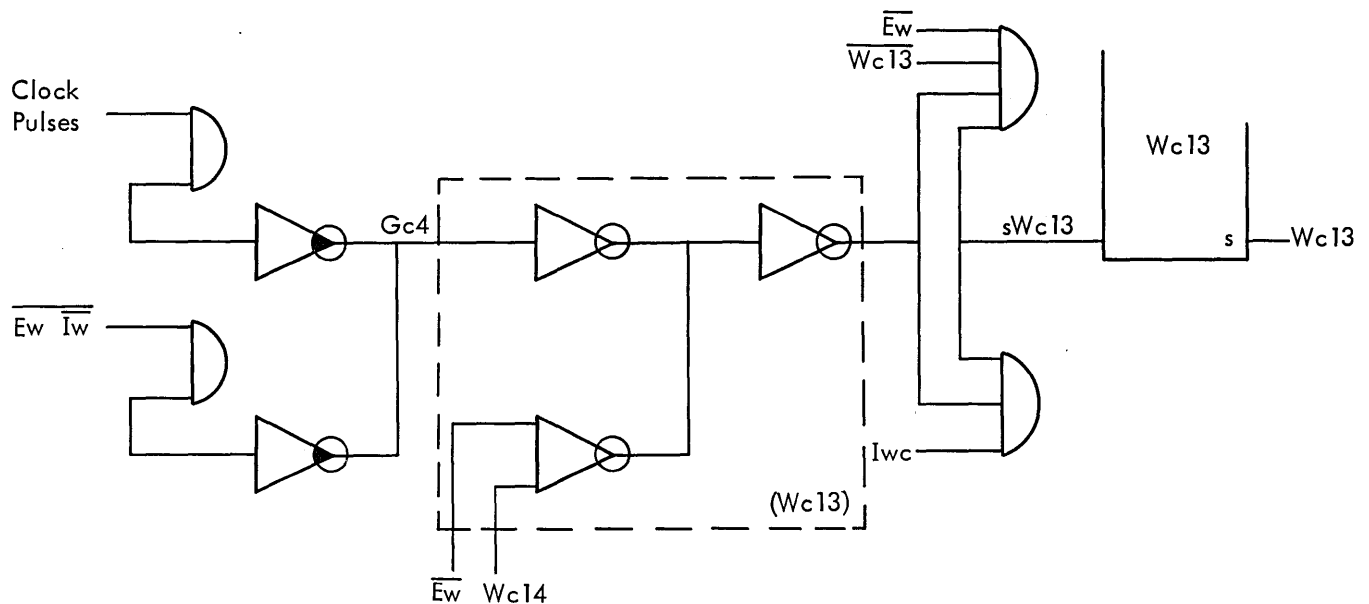


Figure 3-24. Interlace Word Counter - Typical Clock Input

3. 108 The Address Counter is set up in much the same way as the Word Counter. Again the counter counts up, but in this case the one's complement is not used. With the Address Counter, the actual address of the first memory location to be used is placed in the register. As with the Word Counter, Iwc clears the counter then Iwp or Iwe sets up each stage with data from the C-Register. For some stages of the counter, the flip-flops are reset by Iwc then set with the data. For other stages, the flip-flops are set by Iwc then reset with the complement of the data. Either way, the same thing is accomplished. It was a matter of convenience in wiring for one method or the other to have been used.

3. 109 The counting of the Address Counter is handled in the same way as the Word Counter. The least significant stages of both counters are triggered by $Iwa \overline{Ew}$.

$$Iwa = R_{xw} T_{sm}$$

R_{xw} is the Time Share Select Flip-Flop and is covered in detail in paragraph 3. 113. T_{sm} is a signal from the CPU indicating that counter information for each word has been received by the CPU.

3. 110 In the instruction sequence, EOM, EOM, POT, used to set up and start the interlace operation, the second EOM may be omitted if the most significant bits of the Word and Address Counters are "zeros". Refer to figure 3-25 for the relationship between the instruc-

tion bits stored in the C-Register and their respective positions in the counters.

3. 111 The POT instruction which resets the Interlace Prepare Flip-Flop, Ew, also sets the Interlace Active Flip-Flop, Iw.

$$sIw = Pot \ 1 \ (T_3 - T_0) \ Ew \ \overline{Iw}$$

Ew furnishes the ready signal required for the POT instruction. Both the Ew and Iw flip-flops inhibit W-channel interrupts.

$$IIw = \overline{Wf} \ W0 \ \overline{Wh} \ (En + \textcircled{En}) \\ \overline{Iw} \ \overline{Ew} \ \overline{Iwg}$$

3. 112 Refer to figure 3-26 for the timing for the interlace loading process. The Interlace Active Flip-Flop, Iw, allows the buffer to issue a time share request (Trqw) to the CPU whenever the channel needs access to memory.

$$Trqw = \overline{Wf} \ W0 \ \overline{Wh} \ Iw \ \overline{Iwf}$$

Trqw is combined with similar signals from the other channels to produce a common request term for all TMCC's.

$$Trqx = Trq(c) + Trq(d)$$

$$Trq = Trqx + Trqw + Trqy$$

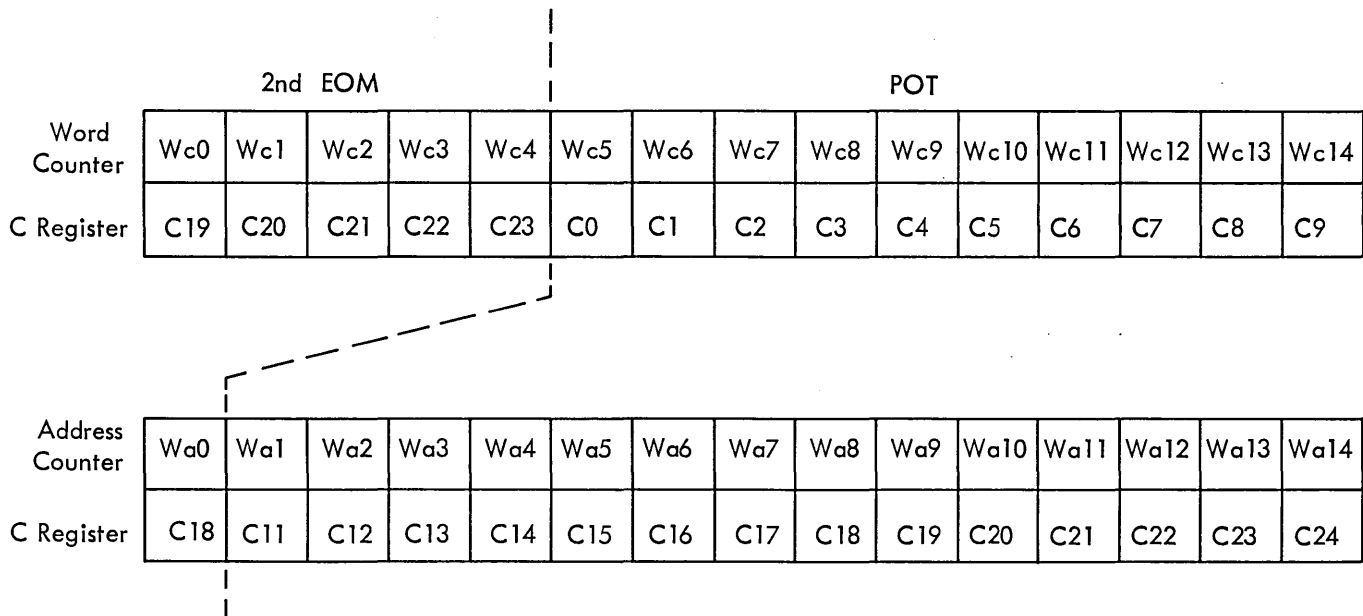


Figure 3-25. Relationship of Instruction Bits to Address and Word Counter Bits

3. 113 When two or more TMCC's make a time share request at the same time or when one channel makes a request while another is already performing a time share operation, the TMCC priority logic determines which channel is allowed to access memory first. For this purpose, each channel has its own Time Share Select flip-flop. No more than one of these flip-flops may be set at any time. The priority established by the flip-flops for the four TMCC channels is D, C, Y, and W in descending order. However, once a channel has been selected, it is allowed to complete the processing of its word without being disturbed by other channels. Using the Time Share Request signals mentioned previously, the priority logic to select the channel is:

Channel W; $sRwx = \overline{Tsm} Trqw \overline{Trqy} \overline{Trqx} (T7 - T0)$
 $rRwx = \overline{Tsm} T0$

Channel Y; $sRyx = \overline{Tsm} Trqy \overline{Trqx} (T7 - T0)$
 $rRyx = \overline{Tsm} T0$

3. 114 The TMCC's for the C and D channels are identical to those for the W and Y channels and so use the same logic nomenclature as above. But a request signal, $Trqx$, is not brought into the C and D units from other higher priority TMCC's. Then, for simplicity the C and D channel equations may be written:

Channel C, $sRcx = \overline{Tsm} Trq (c) \overline{Trq (d)} (\overline{T7} - T0)$
 $rRcx = \overline{Tsm} T0$

Channel D, $sRdx = \overline{Tsm} Trq (d)$
 $rRdx = \overline{Tsm} T0$

The terms Rcx and Rdx are used here for clarity but do not actually appear in the logic equations for the equipment.

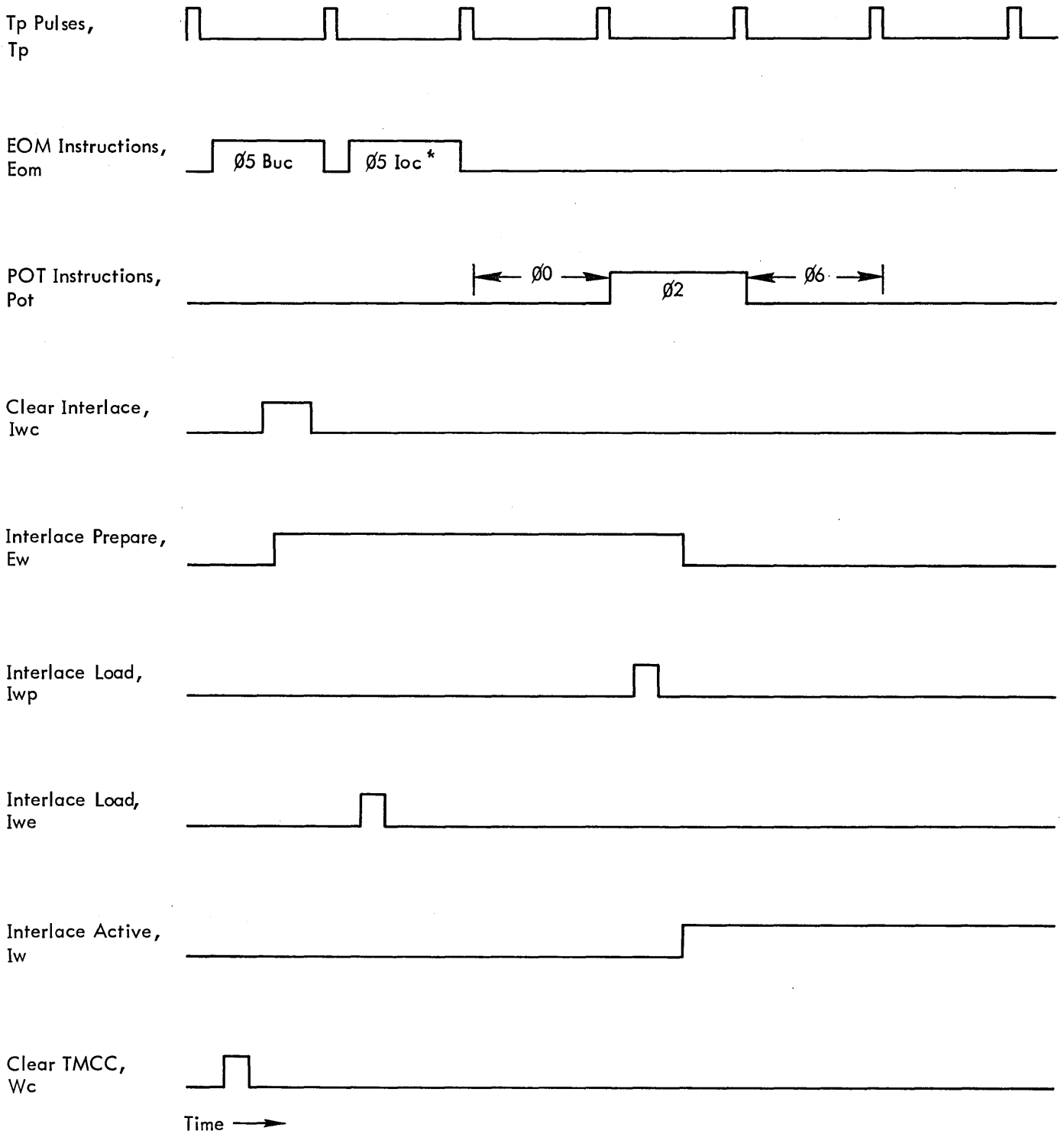
3. 115 If a TMCC channel makes a time share request and has the highest priority of those making such a request, and if \overline{Tsm} is true, then the channel can set its Time Share Select flip-flop. \overline{Tsm} indicates that the CPU is not already engaged in a time share operation. Confining the discussion to the W channel, Rwx would set at pulse time $Tr - T8$.

3. 116 The Interlace Prepare flip-flop (Ew) also produces Er which prevents output signals, Eom and Ioc , from reaching external devices while initializing the Interlace.

$$\overline{Er} = \overline{Ew} \overline{Ey} \dots$$

$$Ioc = Ioc \ 1 \ \overline{CT} \ \overline{Er} \ Qr3$$

$$\textcircled{Eom} = \overline{Eom} \ \overline{Er}$$



*Second EOM may be eliminated if not needed to set count in most significant bits of Address and Word Counters or to set up the extended mode.

Figure 3-26. Interlace Register Loading Timing Chart

The CPU then answers the request with a Tsr signal which remains on for two machine cycles. Tsr enables Wxx to permit a data transfer between W and C registers. While Tsr is true, the CPU sets Tsm to increment the Word and Address Counters.

$$Iwa = Rwx \overline{Tsm} = \text{count trigger}$$

$$W_{xx} = Rwx \overline{Tsr} \overline{Iw} = \text{load buffer from C}$$

3. 117 Refer to figure 3-27 for the timing of the signals involved in the data transfer between the W and C Registers. Figure 3-27, indicates that Wxx remains on for two machine cycles. This allows the data exchange to take place twice. The two cycles are used as follows.

a. Input Operation. The first cycle exchanges the input word and any word currently in the C-Register. Then between cycles, the input word is copied from the C-Register into memory (in parallel). During the second exchange cycle, the word temporarily stored in the W-Register is returned to the C-Register where it can continue in whatever function it may have been participating.

b. Output Operation. The first cycle shifts any information which was being operated on in the C-Register to the W-Register for temporary storage. The word requested by the TMCC Interlace Address Counter is then parallel transferred from memory to the C-Register (between cycles). The second exchange cycle then returns the word from the W-Register back to the C-Register and brings the word out of memory from the C to the W-Register.

3. 118 The double exchange cycle thus provides the TMCC with a means of:

a. moving a word from the W-Register through the C-Register to memory or vice versa, and

b. preserving the contents of the C-Register while the transfer takes place. This is important if the CPU is engaged in some form of computation when the time share takes place. Figure 3-28 is an illustration of the timing involved for a typical interlaced I/O process.

3. 119 Another important aspect of the W and C-Register data exchange during an interlaced output concerns the path taken by the output word in getting to the W-Register. If a clock has read the last word out of the Character Register when the exchange takes place, the next word is precessed simultaneously with its transfer to the W-Register. This occurs as follows. During both of the exchange cycles, Wxx is on because of Rxw and Tsr.

$$W_{xx} = R_{wx} \overline{Tsr} \overline{Iw} + \dots$$

3. 120 If at this time a clock has already been detected (to read the last character of the previous word) so that W5 is set, W4 is set at pulse time T8 in the beginning of the second exchange cycle.

$$sW4 = W5 \overline{Wf} \overline{T8} \overline{Wg}$$

3. 121 Thus, during the second cycle, Wxx and W4 are both true. The output word may then be shifted directly from the C-Register to the Character Register and on through to the W-Register. A one cycle precession is thereby automatically accomplished and the character may be read by the next clock without waiting for another precession to take place.

$$sRw1 = W4 \overline{Wxx} \overline{C21r} + \dots$$

$$rRw1 = W4 \overline{Wxx} \overline{C21r} + \dots$$

$$sRw2 = W4 \overline{Wxx} \overline{C22r} + \dots$$

$$rRw2 = W4 \overline{Wxx} \overline{C22r} + \dots$$

$$sRw3 = W4 \overline{Wxx} \overline{C23r} + \dots$$

$$rRw3 = W4 \overline{Wxx} \overline{C23r} + \dots$$

$$\left. \begin{aligned} Ww1 &= W4 \overline{Wb1} (T7 - T0) + \dots \\ Ww2 &= W4 \overline{Wb2} (T7 - T0) + \dots \\ Ww3 &= W4 \overline{Wb3} + \dots \end{aligned} \right\} \text{Refer to figure 3-3 for } Wb1, Wb2, \text{ and } Wb3.$$

3. 122 This same path for loading the W-Register through the Character Register is also followed when the interlace is initially set up and the first word is called for, if the EOM instruction calls for starting without leader. In this case W5 is first set by Ws C13 instead of a clock signal from the peripheral unit.

$$sW5 = Ws \overline{C13} \overline{C18} \overline{W10} \overline{W11} \overline{W12} \overline{W13} \overline{W14}$$

Figure 3-29 illustrates the use of this feature in loading the first word.

3. 123 When transferring data directly from the C to the Character Register, the parity bit is generated as previously described except that Cpr is used in place of (Wn1 ⊕ Wn2 ⊕ Wn3). Cpr performs a similar function but originates in the CPU and monitors the parity of the octal groups coming from the C-Register rather than from the W-Register.

$$sRwp = W9 \overline{W4} \overline{Rwp} \overline{Wxx} \overline{Cpr} \overline{Qw2} (T7 - T0) + \dots$$

$$rRwp = W9 \overline{W4} \overline{Rwp} \overline{Wxx} \overline{Cpr} \overline{Qw2} (T7 - T0) + \dots$$

Refer to figure 3-4 for Qw2.

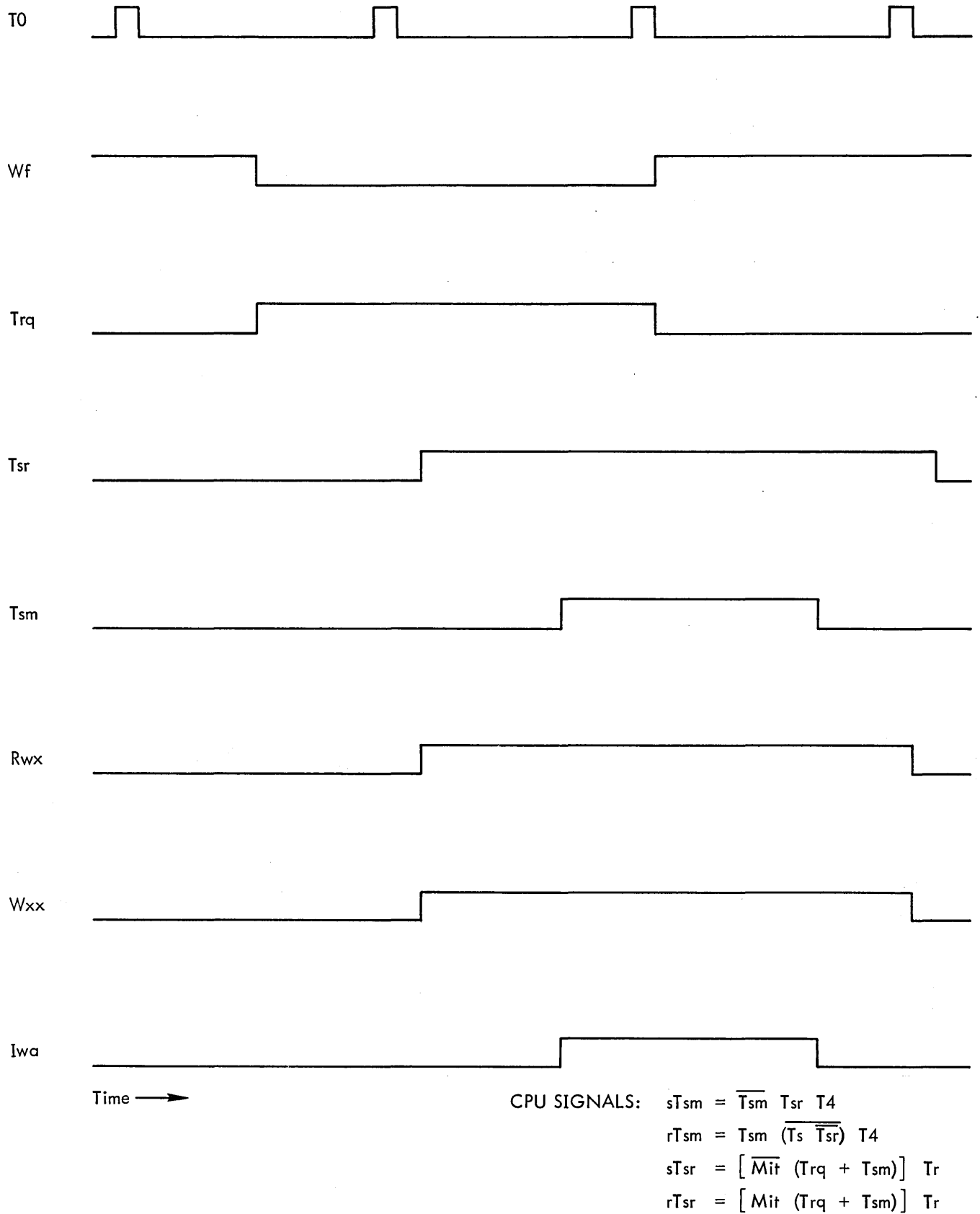


Figure 3-27. Interlace Word Transfer Timing Chart

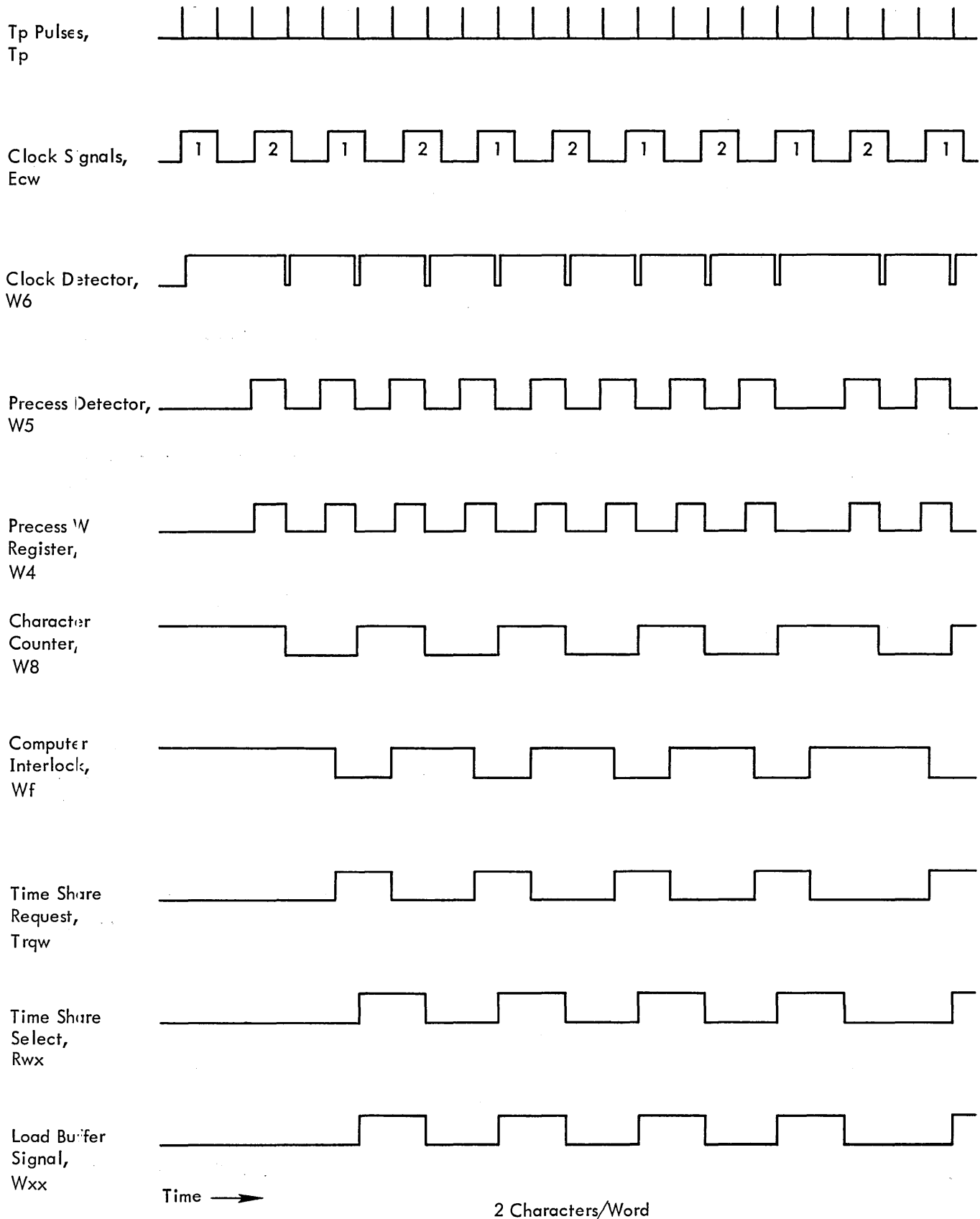


Figure 3-28. Interlace Input/Output Timing Chart

Output without Leader

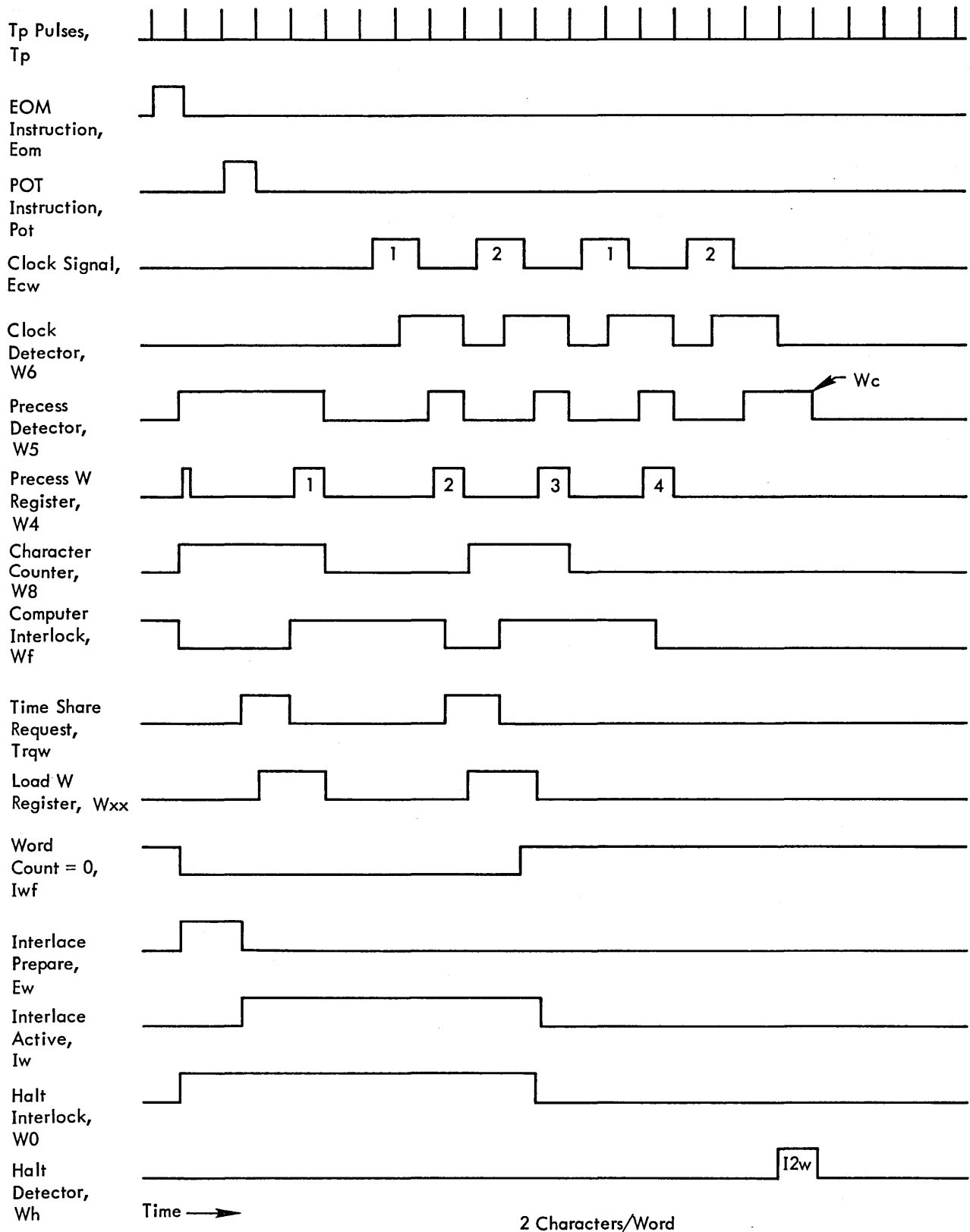


Figure 3-29. Interlace Output Timing Chart

3. 124 As each I/O word is processed, the Word and Address Counters are incremented by Iwa. Iwa also sets Wf again after it has issued the time share request via Trqw.

$$Iwa = Rwx \ Tsm$$

$$sWf = Iwa \ \overline{Mit} \ T0$$

The term \overline{Mit} , when false, denotes that a direct access I/O channel (DACC) is accessing the memory. Since the DACC may require access during an I/O operation by the TMCC, \overline{Mit} is used to momentarily stop the TMCC's action. Memory access by the DACC, although producing Tsm, thus cannot set Wf.

3. 125 When the Word Counter reaches the all "ones" condition, the count is decoded by Iwf.

$$Iwf = Wc0 \ Wc1 \ Wc2 \ Wc3 \ Wc4 \ Wc5 \ Wc6 \ Wc7$$

$$Wc8 \ Wc9 \ Wc10 \ Wc11 \ Wc12 \ Wc13 \ Wc14 \ \overline{Ew}$$

Iwf inhibits further time share request signals and resets the Interlace Active flip-flop.

$$Trqw = \overline{Wf} \ \overline{W0} \ \overline{Wh} \ Iw \ \overline{Iwf} = \text{time share request}$$

$$rIw = Iwf \ T8 = \text{Interlace Active}$$

3. 126 If the Interlace is controlling an input process, the next word loaded into the W-Register after Iw is reset, generates a Word Ready Interrupt, I1w.

$$I1w = \overline{Wf} \ \overline{W0} \ \overline{Wh} \ (En + \overline{En}) \ \overline{Iw} \ \overline{Ew} \ \overline{Iwg}$$

At this time the program can reload the interlace if reading is to continue. On channel W (or Y) the contents of the buffer can be stored with a WIM (or YIM) instruction.

3. 127 During the input process, if an End-of-Record is encountered before the word count is completed, termination takes place as described earlier for the non-interlaced input. Wg detects the gap and sets Wh.

$$sWh = Wg \ \overline{Iwg} \ T8$$

An interrupt is generated and the TMCC is cleared in the usual manner.

$$I2w = (En + \overline{En}) \ \overline{Iwg} \ Wh \ \overline{Wf} + \dots$$

3. 128 If the Interlace is controlling an output process, Iwf resets the Halt Interlock flip-flop, W0, and blocks I1w and Trqw.

$$rW0 = W9 \ Iw \ Iwf \ (\overline{Iwg} + \dots) \ (\overline{T7} - T0)$$

$$I1w = \overline{Wf} \ \overline{W0} \ \overline{Wh} \ (En + \overline{En}) \ \overline{Iw} \ \overline{Ew} \ \overline{Iwg}$$

$$Trqw = \overline{Wf} \ \overline{W0} \ \overline{Wh} \ Iw \ \overline{Iwf}$$

Thus, Rwx is not set again and Iwa cannot be turned on.

$$sRwx = \overline{Tsm} \ Trqw \ \overline{Trqy} \ \overline{Trqx} \ (\overline{T7} - T0)$$

$$Iwa = Rwx \ Tsm$$

Then when the last word is processed out of the W-Register, Wf is not set again.

$$sWf = Iwa \ \overline{Mit} \ T0$$

This results in a situation similar to that for the non-interlaced output. The condition, $\overline{W0} \ \overline{W4} \ \overline{W5} \ \overline{W6}$, exists following the precession of the last character. The state, $\overline{W0} \ \overline{W5} \ \overline{W6}$, then sets Wh. Or, if magnetic tape is being used, $\overline{W0} \ \overline{W5} \ \overline{W6}$ is detected by the tape unit and after a delay, Whs is generated to set Wh.

$$sWh = W9 \ \overline{W11} \ \overline{W0} \ \overline{W5} \ \overline{W6} \ (\overline{Iwg} + \dots)$$

$$T8 + Whs \ W11 \ T8 + \dots$$

Then, halt interrupt and clear signals are generated and the TMCC, including Wh, is reset.

$$I2w = (En + \overline{En}) \ \overline{Iwg} \ Wh \ \overline{Wf} + \dots$$

$$Wc = Wh \ \overline{Wf} \ (T3 - T0) + \dots$$

$$rWh = Wh \ \overline{Wf} \ T8 + \dots$$

3. 129 The termination timing for a typical interlaced input process is illustrated in figure 3-30. In general, termination of a Compatible Mode I/O process is much the same as that for a non-interlaced I/O operation, but figure 3-30 illustrates the relationship of the additional interlace signals involved with those shown in earlier figures.

3. 130 As with other forms of I/O operation, a disconnect EOM (address 00) can also terminate an interlaced operation. This is done by resetting the Interlace Active Flip-Flop and the Extend Operations Flip-Flop through Ws and Wsc. (Refer to the paragraphs on Extended Mode beginning with 3-100 for explanation of Iwg.)

$$rIw = Ws \ \overline{C19} \ \overline{C20} \ \overline{C21} \ \overline{C22} \ \overline{C23} \ (T3 - T0) \ Iw + \dots$$

$$rIwg = Wsc \ T8 \ \overline{W10} \ \overline{W11} \ \overline{W12} \ \overline{W13} \ \overline{W14} \ Iwg + \dots$$

$$sWsc = Ws \ \overline{C19} \ \overline{C20} \ \overline{C21} \ \overline{C22} \ \overline{C23} + \dots$$

3. 131 INTERLACE, EXTENDED MODE (Time Share)

3. 132 The Word Counter and Address Counter set-up procedures are similar for both the Compatible and Extended

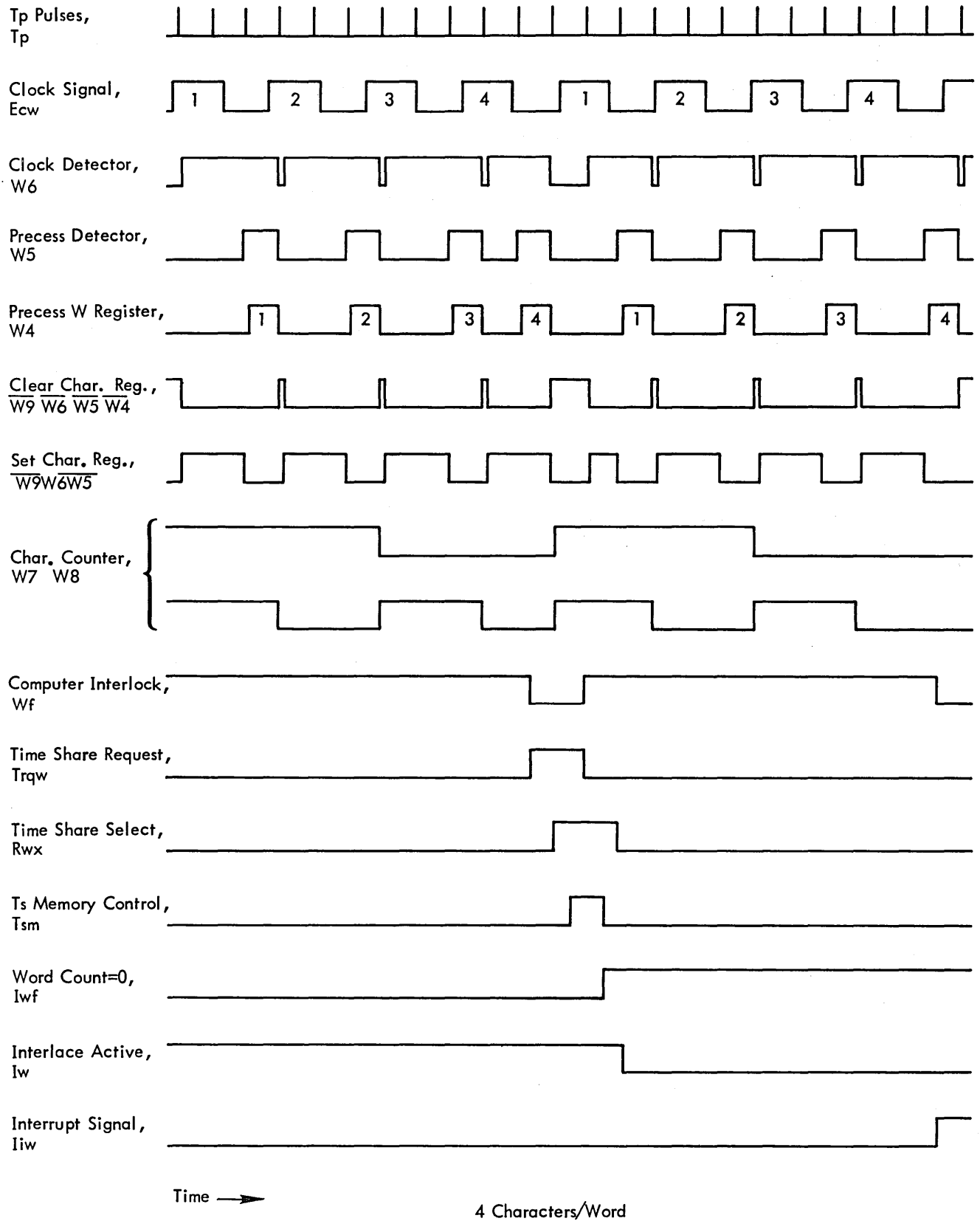


Figure 3-30. Input Termination Timing Chart - Interlace (Compatible Mode)

Modes. That is, the same sequence of instructions is used, (EOM-EOM-POT). However, to select the Extended Mode of operation, it is necessary to place a "one" in bit 12 of the second EOM (Ioc) instruction. When this instruction is processed, the Iwe loading signal sets bit 12 into the Extended Operations Flip-Flop, Iwg.

$$Iwe = Ioc \ 1 \ (T6 + T5) \ Ew$$

$$sIwg = Iwe \ C12$$

Prior to this, Iwg was cleared by the first EOM instruction as was the rest of the Interlace logic.

$$Iwc = Eom \ C9 \ \overline{C10} \ \overline{C1} \ \overline{C17} \ (T3 - T0) \ \overline{C1} \ \text{becomes } C1 \ \text{for TMCC-C}$$

$$rIwg = Iwc + . . .$$

3. 133 To use the Extended Mode, four additional flip-flops are loaded by the Ioc instruction. Two of these, Iwh and Iwi, comprise the Channel Command Register which selects the type of termination. The remaining two flip-flops, Iwj and Iwk, are employed to arm the I2w and I1w interrupts on a selective basis.

$$sIwj = Iwe \ C13$$

$$rIwj = Iwc + . . .$$

$$sIwk = Iwe \ C14$$

$$rIwk = Iwc + . . .$$

$$sIwh = Iwe \ C15$$

$$rIwh = Iwc$$

$$sIwi = Iwe \ C16$$

$$rIwi = Iwc$$

By decoding Iwj and Iwk, the Extended Mode can perform the four different terminal functions listed in table 3-4. Each function can be used to control either an input or an output operation. In the following paragraphs, each of the four functions is discussed in detail. In each case, it is assumed that the interlace registers have already been loaded by the EOM-EOM-POT instructions and that the interlaced I/O operation is proceeding normally.

3. 134 IORD - Iwg \overline{Iwh} \overline{Iwi}

3. 135 Output

3. 136 Write C words. When C equals zero, output is terminated (i.e. the device is signaled that the last characters have been transmitted). When the peripheral device has generated the End-of-Record and, if

necessary, checked the validity of the record, it sends an End-of-Record response to the channel buffer. When received by the buffer, the End-of-Record signal generates an End-of-Record interrupt (if armed) and disconnects the channel.

3. 137 The line printer generates the End-of-Record response when it completes the printing of a line. If the printer encounters any print errors or faults, it sends a signal to the channel that sets the channel error indicator. This can occur since the printer has not disconnected from the channel. The IORD is useful when the program is to print several lines and the program is not otherwise to use the channel between lines. When the printer completes each line, it causes an End-of-Record interrupt (assumed to be armed), notifying the program that it can immediately transmit the next paper control instruction and the next line image.

3. 138 The unbuffered card punch operates similarly. It generates the End-of-Record response after punching each row. If any faults occur during the punching of the entire card, the card punch sends a signal to the channel that sets the channel error indicator; this occurs after punching the last row (row 9).

NOTE

A program should not use IORD with devices that do not have End-of-Record conditions on output (e.g., devices such as the paper tape punch and typewriter). These devices to terminate output but give the program no indication when they receive the last characters.

After the last word is accessed from memory, zero word count is established.

3. 139 The interlace is counted by the Interlace Count Trigger, Iwa.

$$Iwa = Rwx \ Tsm$$

Zero Word Count, Iwf, occurs and the Interlace Active flip-flop, Iw is reset.

$$Iwf = Wc0 \ Wc1 \ . . . \ Wc13 \ Wc14 \ \overline{Ew}$$

$$rIw = Iwf \ T8$$

The Halt Interlock, W0, is reset.

$$rW0 = W9 \ Iw \ Iwf \ (\overline{Iwg} + \overline{Iwh} + \overline{Iwi}) \ (\overline{T7 - T0}) + . . .$$

If the End-of-Transmission Interrupt Enable, Iwk, has been previously armed, an I1w interrupt occurs and Iwk is reset.

$$I1w = Iwg \ \overline{Iw} \ Iwf \ Iwk + . . .$$

Table 3-4. Interlace Extended - Mode Terminal Functions

Terminal Function	Iwg	Iwh	Iwi	Summary of Operation
IORD Input/Output of record then disconnect.	1	0	0	The I/O operation proceeds until the word count equals zero then terminates. On input, the channel disconnects when the End-of-Record is encountered. On output, the channel signals the device that the last character has been transmitted then disconnects after the device provides an End-of-Record response.
IOSD Input/Output until signal then disconnect.	1	0	1	The channel disconnects when the word count equals zero or at the end of a record.
IORP Input/Output of a record then proceed.	1	1	0	The I/O operation proceeds until the word count equals zero but does not terminate. On input, the channel sets the inter-record indicator when the end of a record is encountered. On output, the channel signals the device that the last character has been transmitted then sets the inter-record indicator after the device provides an End-of-Record response. The channel does not disconnect (except for magnetic tape).
IOSP Input/Output until signal then proceed.	1	1	1	When the word count equals zero, the program should either reload the interlace to continue, or terminate the operation before the next clock is received; otherwise a rate error will occur.

$$rIwk = I1w Iwk T8 + \dots$$

When the last character of the last word has been pre-processed into the Character Register, Wf is reset.

$$rWf = \overline{W7} \overline{W8} W4 (T6 + T5) + \dots$$

Because the Halt Interlock, W0, and Interlace Active, Iw, have been reset, a Time Share Request, Trqw, can not be made.

$$Trqw = \overline{Wf} W0 \overline{Wh} Iw \overline{Iwf}$$

Any further clocking, Ecw, of the buffer generates a Halt Interlock Signal condition.

$$sW6 = \overline{W5} Ecw T8 \overline{W10} \overline{W11} \overline{W12} \overline{W13} \overline{W14}$$

$$sW5 = \overline{W5} W6 \overline{Ecw} T0 + \dots$$

$$rW6 = W5 T0 + \dots$$

Halt Interlock Signal (decoded by the peripheral device)

$$= W5 \overline{W6} \overline{W0}$$

If the Halt Signal, Whs, is received the Halt Detector Wh, is set and a disconnect occurs.

$$sWh = W9 \overline{Iwh} T8 Whs + \dots$$

$$Wc = Wh \overline{Wf} (T3 - T0) + \dots$$

$$rW9 = Wc$$

$$rW10 = Wc$$

etc.

The Signal Complete flip-flop, Wsc, is set.

$$sWsc = Wh \overline{Wf} T8 + \dots$$

If the End-of-Record Interrupt Enable, Iwj, has been previously armed, an I2w interrupt occurs and Iwj is reset.

$$I2w = Wsc Iwj Iwg + \dots$$

$$rIwj = I2w Iwj T8 + \dots$$

The Extend Operations flip-flop, Iwg, and the Signal Complete flip-flop, Wsc are then reset.

$$rIwg = Wsc T8 \overline{W10} \overline{W11} \overline{W12} \overline{W13} \overline{W14} Iwg + \dots$$

$$rWsc = Wsc \overline{Iwg} T8$$

3. 140 IOSD - $Iwg \overline{Iwh} Iwi$

3. 141 Output

3. 142 Write C words. When C equals zero and when the last character has been transmitted, the channel disconnects the device and becomes inactive. If an End-of-Record signal is received before the count reaches zero, the channel disconnects immediately.

NOTE

The IOSD is designed for use on devices which are normally operated on the basis of the word count only. Typewriters and paper tape devices are of this type, as are the printer and card punch when the user does not wish to stay connected until the operation is complete.

3. 143 The interlace is counted by the Interlace Count Trigger, Iwa .

$$Iwa = Rwx Tsm$$

Zero Word Count, Iwf , occurs and the Interlace Active flip-flop, Iw , is reset.

$$Iwf = Wc0 Wc1 \dots Wc13 Wc14 \overline{Ew}$$

$$rIw = Iwf T8 + \dots$$

The Halt Interlock, $W0$, is reset.

$$rW0 = W9 Iw (\overline{Iwg} + \overline{Iwh} + \overline{Iwi}) (\overline{T7 - T0}) + \dots$$

If the End-of-Transmission Interrupt Enable, Iwk , has been previously armed, an $I1w$ interrupt occurs and Iwk is reset.

$$I1w = Iwg \overline{Iw} Iwk Iwk + \dots$$

$$rIwk = I1w Iwk T8 + \dots$$

When the last character of the last word has been processed into the Character Register, Wf is reset.

$$rWf = \overline{W7} \overline{W8} W4 (T6 + T5) + \dots$$

Because the Halt Interlock, $W0$, has been reset, a Time Share Request, $Trqw$, cannot be made.

$$Trqw = \overline{Wf} \overline{W0} \overline{Wh} Iw \overline{Iwf}$$

Any further clocking, Ecw , of the buffer generates a Halt Interlock Signal condition.

$$sW6 = \overline{W5} Ecw T8 \overline{W10} \overline{W11} \overline{W12} \overline{W13} \overline{W14}$$

$$sW5 = \overline{W5} W6 \overline{Ecw} T0 + \dots$$

$$rW6 = W5 T0 + \dots$$

Halt Interlock Signal

$$= W5 \overline{W6} \overline{W0}$$

For devices other than magnetic tape, the Halt Detector, Wh , sets upon reaching zero Word Count, Iwf , and after the last character has been clocked from the buffer.

$$sWh = W9 \overline{W11} \overline{W0} W5 \overline{W6} (\overline{Iwh} Iwi + \dots) T8 + \dots$$

The Halt Detector, Wh , also sets upon the occurrence of a Halt Signal, Whs , from the magnetic tape unit.

$$sWh = W9 \overline{Iwh} Whs T8 + \dots$$

The setting of the Halt Detector, Wh , initiates a buffer disconnect sequence.

$$Wc = Wh \overline{Wf} (T3 - T0) + \dots$$

$$rW9 = Wc$$

$$rW10 = Wc$$

etc.

The Signal Complete flip-flop, Wsc , is set.

$$sWsc = Wh \overline{Wf} T8$$

If the End-of-Record Interrupt Enable, Iwj , has been previously armed, an $I2w$ interrupt occurs and Iwj is reset.

$$I2w = Wsc Iwj Iwg + \dots$$

$$rIwj = I2w Iwj T8 + \dots$$

The Extend Operations flip-flop, Iwg , and the Signal Complete flip-flop, Wsc , are then reset.

$$rIwg = Wsc T8 \overline{W10} \overline{W11} \overline{W12} \overline{W13} \overline{W14} Iwg + \dots$$

$$rWsc = Wsc \overline{Iwg} T8$$

3. 144 IORP - $Iwg Iwh \overline{Iwi}$

3. 145 Output

3. 146 Write C words. When the channel interlace counts C down to zero, the interlace notifies the channel buffer that it has received the last word that is to be

output; when the buffer outputs this last word, it sends a signal to the connected peripheral device indicating that the device has the last word. When the peripheral device receives, outputs and checks the validity of this last word, it sends an End-of-Record response to the channel buffer. When received by the buffer, the End-of-Record signal generates an End-of-Record interrupt (if armed) and sets the Inter-Record indicator; the channel does not disconnect.

3. 147 When the peripheral device is magnetic tape, the tape continues to move after it signals End-of-Record. As in reading tape, the signal causes the Tape Gap signal to come high. If the program executes a new write tape or erase tape EOM during the inter-gap time (approximately one millisecond), the tape remains in motion and proceeds to write or erase a new record. If the program executes no such EOM before the Tape Gap signal drops, the channel disconnects and tape comes to a stop. No interrupt occurs at this time. This is the only condition which causes a channel to disconnect automatically for an IORP.

3. 148 To proceed after the End-of-Record occurs, the program first executes a Buffer Control mode EOM to re-initialize the Channel Unit Address Register and then reloads the interlace portion of the channel (the program can alert the Interlace via the Buffer Control EOM). Otherwise, the channel immediately terminates any attempt to use its interlace portion, since the channel is still active and in the End-of-Record condition. When the program continues from an Inter-Record condition, the program should use an extended mode terminal function.

NOTE

A program should not use IORP with devices that do not generate End-of-Record responses upon output termination; such devices are paper tape and typewriter. These devices do terminate output but give the program no indication when they receive the last characters. The IORP should also not be used with the printer and card punch since these devices expect the channel to disconnect after they send EOR.

3. 149 After the last word is accessed from memory, the Interlace is counted and zero word count is established. The Interlace is counted by the Interlace Count Trigger, Iwa.

$$Iwa = Rwx Tsm$$

Zero Word Count, Iwf, occurs and the Interlace Active flip-flop, Iw, is reset.

$$Iwf = Wc0 Wc1 \dots Wc13 Wc14 \overline{Ew}$$

$$rIw = Iwf T8 + \dots$$

The Halt Interlock, W0, is reset.

$$rW0 = W9 Iwf (\overline{Iwg} + \overline{Iwh} + \overline{Iwi}) \\ (\overline{T7 - T0}) + \dots$$

If the End-of-Transmission Interrupt Enable, Iwk, has been previously armed, an I1w interrupt occurs and Iwk is reset.

$$I1w = Iwg \overline{Iw} Iwf Iwk + \dots$$

$$rIwk = I1w Iwk T8 + \dots$$

When the last character of the last word has been processed into the Character Register, Wf is reset.

$$rWf = \overline{W7} \overline{W8} W4 (T6 + T5) + \dots$$

Because the Halt Interlock, W0, and Interlace Active, Iw, have been reset, a Time Share Request, Trqw, cannot be made.

$$Trqw = \overline{Wf} W0 \overline{Wh} Iw \overline{Iwf}$$

Any further clocking, Ecw, of the buffer generates a Halt Interlock Signal condition.

$$sW6 = \overline{W5} Ecw T8 \overline{W10} \overline{W11} \overline{W12} \overline{W13} \overline{W14}$$

$$sW5 = \overline{W5} W6 \overline{Ecw} T0 + \dots$$

$$rW6 = W5 T0 + \dots$$

Halt Interlock Signal

$$= W5 \overline{W6} \overline{W0}$$

3. 150 For non-magnetic tape devices, the buffer awaits the receipt of a Halt Signal, Whs, from the device. The End-of-Record Detector, Wg, is set.

$$sWg = Whs (\overline{T7 - T0}) \overline{W10} \overline{W11} \overline{W12} \overline{W13} \overline{W14} + \dots$$

The Signal Complete Detector, Wsc, sets and if the End-of-Record Interrupt Enable, Iwj, has been previously armed, an I2w interrupt is generated and Iwj is reset.

$$sWsc = Wg Iwg Iwh T8 (Iwf + \dots) + \dots$$

$$I2w = Wsc Iwj Iwg + \dots$$

$$rIwj = I2w Iwj T8 + \dots$$

The device does not disconnect because Wh is not permitted to set.

3. 151 For magnetic tape, the buffer awaits the receipt of a gap signal, Mtgw. The End-of-Record Detector, Wg, is set.

$$sWg = Mtgw T0 Iwg W11 (\overline{W0} \overline{W5} \overline{W6} \overline{W9} + \dots) + \dots$$

The Signal Complete Detector, Wsc, sets and if the End-of-Record Interrupt Enable, Iwj, has been previously armed, an I2w interrupt is generated and Iwj is reset.

$$sWsc = Wg Iwg Iwh T8 (Iwf + \dots) + \dots$$

$$I2w = Wsc Iwj Iwg$$

$$rIwj = I2w Iwj T8 + \dots$$

The magnetic tape system can continue if the interrupt sub-routine executes an EOM to the tape within approximately one millisecond from the occurrence of the interrupt. If no EOM is executed, the tape generates a Halt Signal, Whs, and the Halt Detector, Wh, sets. The magnetic tape is disconnected and the buffer is cleared.

$$sWh = Whs W11 T8 + \dots$$

$$Wc = Wh \overline{Wf} (T3 - T0) + \dots$$

$$rW9 = Wc$$

$$rW10 = Wc$$

etc.

$$rIwg = Wsc T8 \overline{W10} \overline{W11} \overline{W12} \overline{W13} \overline{W14} Iwg + \dots$$

$$rWsc = Wsc T8 \overline{Iwg}$$

$$rWh = Wh \overline{Wf} T8 + \dots$$

3. 152 IOSP - Iwg Iwh Iwi

3. 153 Output

3. 154 Write C words. When the channel counts C down to zero, the channel generates a Count Equals Zero interrupt (if armed); the channel does not terminate output. The program should reload the interlace portion of the channel to continue writing in the same record. Failure to reload the Interlace before the buffer transmits all of the characters in its registers and before the peripheral device requests the next character from

the buffer results in a rate error; this sets the Channel Error Indicator.

3. 155 If the program executes a TERMINATE OUTPUT (TOP) instruction after the channel has counted C down to zero, the channel terminates the output and operates identically like the IORP from this point on.

3. 156 After the last word is accessed from memory, the interlace is counted and zero word count is established.

The Interlace is counted by the Interlace Count Trigger, Iwa.

$$Iwa = Rwx Tsm$$

Zero Word Count, Iwf, occurs and the Interlace Active flip-flop, Iw, is reset.

$$Iwf = Wc0 Wc1 \dots Wc13 Wc14 \overline{Ew}$$

$$rIw = Iwf T8 + \dots$$

If the End-of-Transmission Interrupt Enable, Iwk, has been previously armed, an I1w interrupt occurs and Iwk is reset.

$$I1w = Iwg \overline{Iw} Iwf Iwk + \dots$$

$$rIwk = I1w Iwk T8 + \dots$$

The program, upon receipt of an I1w interrupt, should reload the Interlace to permit the transmission to continue with a new set of parameters (i. e., word count, address, terminal functions, etc.). The loading of the Interlace with a zero word count could permit conversion of the current IOSP to some other terminal function, for example to an IORD, thereby effecting a disconnect.

3. 157 Execution of a TERMINATE OUTPUT (TOP) instruction, e. g., EOM14000, would convert the IOSP to an IORP. The TOP instruction resets the Halt Interlock flip-flop, W0.

$$rW0 = Ioc C12 \overline{C17} \overline{C19} \overline{C20} \overline{C21} \overline{C22} \overline{C23} W9 T0 + \dots$$

Any further clocking of the channel by Ecw generates a Halt Interlock Signal condition when the last character has been clocked from the buffer.

$$sW6 = \overline{W5} Ecw T8 \overline{W10} \overline{W11} \overline{W12} \overline{W13} \overline{W14}$$

$$sW5 = \overline{W5} W6 \overline{Ecw} T0 + \dots$$

$$rW6 = W5 T0 + \dots$$

$$rW5 = W4 T0 + \dots$$

Halt Interlock Signal

$$= W5 \overline{W6} \overline{W0}$$

The Halt Interlock Signal is representative of the buffer status had the IOSP been an IORP.

3.158 IORD - $Iwg \overline{Iwh} \overline{Iwi}$

3.159 Input

3.160 Read C words. If C equals zero before the End-of-Record is detected, the rest of the record is ignored. At the End-of-Record, the peripheral device is disconnected and the channel becomes inactive.

3.161 When the W-Register acquires the specified number of characters, a Time Share Request, $Trqw$, is generated.

$$Trqw = \overline{Wf} \overline{W0} \overline{Wh} \overline{Iw} \overline{Iwf}$$

When memory is accessed, the interlace registers are counted by the Interlace Count Trigger, Iwa .

$$Iwa = Rwx Tsm$$

Zero Word Count, Iwf , may occur and the Interlace Active flip-flop, Iw , is reset.

$$Iwf = Wc0 Wc1 \dots Wc13 Wc14 \overline{Ew}$$

$$rIw = Iwf T8 + \dots$$

If the Zero Word Count Interrupt Enable, Iwk , has been previously armed, then an Ilw interrupt occurs.

$$Ilw = Iwg \overline{Iw} Iwf Iwk + \dots$$

$$rIwk = Ilw Iwk T8 + \dots$$

Additional characters entering the channel after Zero Word Count has been reached are precessed into the W-Register.

$$sW6 = \overline{W5} Ecw T8 \overline{W10} \overline{W11} \overline{W12} \overline{W13} \overline{W14}$$

$$rW6 = W5 T0 + \dots$$

$$sW5 = \overline{W5} \overline{W6} \overline{Ecw} T0 + \dots$$

$$rW5 = W4 T0 + \dots$$

$$sW4 = W5 Wf T8 \overline{Wg} + \\ + Iwg \overline{W9} \overline{Iwi} W5 Iwf T8 + \dots$$

$$sW4 = Wr T0 + \dots$$

$$Ww1 = W4 Wb1 (T7 - T0) + \dots$$

$$Ww2 = W4 Wb2 (T7 - T0) + \dots$$

$$Ww3 = W4 Wb3 + \dots$$

However, Time Share Request, $Trqw$, is inhibited.

$$Trqw = \overline{Wf} \overline{W0} \overline{Wh} \overline{Iw} \overline{Iwf}$$

Parity errors cannot occur after Zero Word Count.

$$sWe = \overline{W9} \overline{W6} \overline{W5} \overline{W4} Rwp \overline{Wg} Npw \\ (\overline{Iwg} + \overline{Iwi} + \overline{Iwf}) + \dots$$

Rate errors cannot set We while $W4$ is enabled by Iwg $W9$ Iwi Iw .

$$sWe = W0 \overline{W6} W5 Ecw T8 + \dots$$

After Zero Word Count is established, detection of a Halt Signal or Photoreader Gap sets the End-of-Record Detector, Wg .

$$sWg = Whs (T7 - T0) \overline{W10} \overline{W11} \overline{W12} \overline{W13} \overline{W14} \\ + \overline{W9} \overline{W10} \overline{W11} \overline{W12} \overline{W13} \\ (\overline{Rw1} \overline{Rw2} \overline{Rw3} \overline{Rw4} \overline{Rw5} \overline{Rw6} \overline{Rwp}) \\ W5 (T7 - T0) \overline{W10} \overline{W11} \overline{W12} \overline{W13} \overline{W14} + \dots$$

The Halt Detector, Wh , sets.

$$sWh = Whs W11 T8 + Wg \overline{W9} \overline{W11} \overline{Iwh} Iwg T8 \\ (Iwf + \dots) + \dots$$

The buffer is flushed (i.e., allowed to precess without receiving input clocks until the Character Count equals zero) until it is assured that Wf is reset.

$$sW4 = Wh Wf T8 + \dots$$

$$rWf = \overline{W7} \overline{W8} W4 (T6 + T5) + \dots$$

The buffer is cleared and the peripheral device disconnected.

$$Wc = Wh Wf (T3 - T0) + \dots$$

If the End-of-Record Interrupt Enable, Iwj , has been previously armed, and End-of-Record Interrupt, $Iw2$, occurs.

$$sWsc = Wh \overline{Wf} T8 + \dots$$

$$rWsc = Wsc T8 \overline{Iwg}$$

$$I2w = Wsc Iwj Iwg + \dots$$

$$rIwj = I2w Iwj T0 + \dots$$

$$rIwg = Wsc Iwg \overline{W10} \overline{W11} \overline{W12} \overline{W13} \overline{W14} T8 + \dots$$

Should an End-of-Record occur before Zero Word Count is established, the End-of-Record Detector is set.

$$\begin{aligned} sWg &= Mtgw T0 Iwg W11 (\overline{W0} \overline{W9} + \dots) \\ &+ Whs (\overline{T7} - \overline{T0}) \overline{W10} \overline{W11} \overline{W12} \overline{W13} \overline{W14} \\ &+ \overline{W9} \overline{W10} \overline{W11} \overline{W12} \overline{W13} \\ &(\overline{Rw1} \overline{Rw2} \overline{Rw3} \overline{Rw4} \overline{Rw5} \overline{Rw6} \overline{Rwp}) W5 \\ &(\overline{T7} - \overline{T0}) \overline{W10} \overline{W11} \overline{W12} \overline{W13} \overline{W14} \end{aligned}$$

Any character remaining in the W-Register are flushed by W4 and a Time Share Request, Trqw, is generated.

$$sW4 = Iwg Wg Wf W0 \overline{Wev} Iw T8 \overline{W7} \overline{W9} \overline{W10} \overline{W11} \overline{Wh} + \dots$$

where:

$$\begin{aligned} Wev &= W8 Wn2 W7 Wn1 + \overline{W8} \overline{Wn2} \overline{W7} \overline{Wn1} \\ &+ \overline{W8} \overline{Wn2} \overline{W7} \overline{Wn1} + W8 Wn2 \overline{W7} \overline{Wn1} \end{aligned}$$

$$rWf = \overline{W7} \overline{W8} W4 (T6 + T5) + \dots$$

$$Trqw = \overline{Wf} W0 \overline{Wh} Iw \overline{Iwf}$$

and $\overline{W7} \overline{W9} \overline{W10} \overline{W11} \overline{Wh}$ indicates that a scan operation is not taking place.

3. 162 As a result of one Time Share operation, Wf is set.

$$sWf = Iwa \overline{Mit} T0 + \dots$$

The Character Counter is reloaded with its original count, making Wev true and the flush operation ceases.

3. 163 The Halt Detector, Wh, is now permitted to set. For Magnetic Tape operation this occurs upon receipt of a Halt Signal, Whs.

$$\begin{aligned} sWh &= Whs W11 T8 + Wg \overline{W9} \overline{W11} \overline{Iwh} \\ &(\overline{Iwf} + \overline{Wev} \overline{Wf}) T8 Iwg + \dots \end{aligned}$$

Before the buffer is cleared, the buffer is again flushed but a Time Share Request, Trqw, is inhibited.

$$sW4 = Wh Wf T8 + \dots$$

$$rWf = \overline{W7} \overline{W8} W4 (T6 + T5) + \dots$$

$$Trqw = \overline{Wf} W0 \overline{Wh} Iw \overline{Iwf}$$

The buffer is then cleared.

$$Wc = Wh \overline{Wf} (T3 - T0) + \dots$$

If the End-of-Record Interrupt Enable, Iwj, has been previously armed, an End-of-Record Interrupt, I2w, occurs.

$$sWsc = Wh \overline{Wf} T8 + \dots$$

$$rWsc = Wsc T8 Iwg$$

$$I2w = Wsc Iwj Iwg$$

$$rIwj = I2w Iwj T8 + \dots$$

$$rIwg = Wsc Iwg \overline{W10} \overline{W11} \overline{W12} \overline{W13} \overline{W14} T8 + \dots$$

3. 164 IOSD - Iwg \overline{Iwh} Iwi

3. 165 Input

3. 166 Read C words. When C equals zero or when the End-of-Record is encountered, the device is disconnected and the channel becomes inactive. If the channel disconnects because of a zero count, an EOR interrupt (if armed) is generated in addition to the count-equal-zero interrupt. If both are armed, the count-equal-zero interrupt occurs first.

When the W-Register acquires the specified number of characters, a Time Share Request, Trqw, is generated.

$$Trqw = \overline{Wf} W0 \overline{Wh} Iw \overline{Iwf}$$

When memory is accessed, the interlace registers are counted by the Interlace Count Trigger, Iwa.

$$Iwa = Rwx Tsm$$

Zero Word Count, Iwf, may occur and the Interlace Active flip-flop, Iw, is reset.

$$Iwf = Wc0 Wc1 \dots Wc13 Wc14 \overline{Ew}$$

$$rIw = Iwf T8 + \dots$$

If the zero Word Count Interrupt Enable, Iwk, has been previously armed, then an I1w interrupt occurs.

$$I1w = Iwg \overline{Iw} Iwf Iwk + \dots$$

$$rIwk = I1w Iwk T0 + \dots$$

The Halt Detector is set.

$$sWh = \overline{W9} \overline{W11} Iwg \overline{Iwh} \overline{Iwi} Iwf T8 W0 + \dots$$

The buffer is flushed, but no Time Share Request, Trqw, may be initiated.

$$sW4 = Wh Wf T8 + \dots$$

$$rW4 = W4 T0 + \dots$$

$$\begin{aligned} rWf &= \overline{W7} \overline{W8} T8 + \dots \\ Trqw &= \overline{Wf} \overline{W0} \overline{Wh} Iw \overline{Iwf} \end{aligned}$$

The channel is cleared and the peripheral device disconnected.

$$Wc = Wh \overline{Wf} (T3 - T0) + \dots$$

If the End-of-Record Interrupt Enable, Iwj , has been previously armed, and End-of-Record Interrupt, $I2w$, occurs.

$$\begin{aligned} sWsc &= Wh \overline{Wf} T8 + \dots \\ rWsc &= Wsc \overline{Iwg} T8 \\ I2w &= Wsc Iwj Iwg + \dots \\ rIwj &= I2w Iwj T8 = \dots \\ rIwg &= Wsc Iwg \overline{W10} \overline{W11} \overline{W12} \overline{W13} \overline{W14} \\ &\quad T8 + \dots \end{aligned}$$

Should an End-of-Record occur before Zero Word Count is established, the End-of-Record Detector is set.

$$\begin{aligned} sWg &= Mtgw T0 Iwg W11 (\overline{W0} \overline{W9} + \dots) \\ &\quad + Whs (T7 - T0) \overline{W10} \overline{W11} \overline{W12} \overline{W13} \overline{W14} \\ &\quad + \overline{W9} \overline{W10} \overline{W11} \overline{W12} \overline{W13} \\ &\quad (\overline{Rw1} \overline{Rw2} \overline{Rw3} \overline{Rw4} \overline{Rw5} \overline{Rw6} \overline{Rwp}) \\ &\quad W5 (T7 - T0) \overline{W10} \overline{W11} \overline{W12} \overline{W13} \overline{W14} \end{aligned}$$

Any characters remaining in the W-Register are flushed and a Time Share Request, $Trqw$, is generated.

$$sW4 = \frac{Iwg Wg Wf W0 \overline{Wev} Iw \overline{W7} \overline{W9} \overline{W10} \overline{W11}}{\overline{Wh} T8} + \dots$$

where:

$$\begin{aligned} Wev &= \overline{W8} \overline{Wn2} \overline{W7} \overline{Wn1} + \overline{W8} \overline{Wn2} \overline{W7} \overline{Wn1} \\ &\quad + \overline{W8} \overline{Wn2} \overline{W7} \overline{Wn1} + \overline{W8} \overline{Wn2} \overline{W7} \overline{Wn1} \end{aligned}$$

$$rWf = \overline{W7} \overline{W8} W4 (T6 + T5) + \dots$$

$$Trqw = \overline{Wf} \overline{W0} \overline{Wh} Iw \overline{Iwf}$$

As a result of the Time Share operation, Wf , is set.

$$sWf = Iwa \overline{Mit} T0$$

When the term Wev is true the flush and store operation ceases. The Halt Detector, Wh , is permitted to set

$$\begin{aligned} sWh &= Whs W11 T8 + Wg \overline{W9} \overline{W11} \overline{Iwh} \\ &\quad (Iwf + Wev Wf) T8 Iwg + \dots \end{aligned}$$

Before the buffer is cleared, the buffer is again flushed but a Time Share Request, $Trqw$, is inhibited.

$$sW4 = Wh Wf T8 + \dots$$

$$rWf = \overline{W7} \overline{W8} W4 (T6 + T5) + \dots$$

$$Trqw = \overline{Wf} \overline{W0} \overline{Wh} Iw \overline{Iwf}$$

The buffer is then cleared.

$$Wc = Wh \overline{Wf} (T3 - T0) + \dots$$

If the End-of-Record Interrupt Enable, Iwj , has been previously armed, an End-of-Record Interrupt, $I2w$, occurs.

$$sWsc = Wh \overline{Wf} T8 + \dots$$

$$rWsc = Wsc \overline{Iwg} T8$$

$$I2w = Wsc Iwj Iwg + \dots$$

$$rIwk = I2w Iwk T8 + \dots$$

$$\begin{aligned} rIwg &= Wsc Iwg \overline{W10} \overline{W11} \overline{W12} \overline{W13} \overline{W14} \\ &\quad T8 + \dots \end{aligned}$$

3.167 Should an IOSD mode be used on input with magnetic tape devices, no disconnect will occur at Iwf . Should additional characters enter the buffer after Iwf , a rate error may occur.

$$sWe = W0 \overline{W6} W5 Ecw T8 + \dots$$

Disconnect occurs upon receipt of a Halt Signal, Whs , from the magnetic tape system.

$$sWh = Whs W11 T8 + \dots$$

$$sW4 = Wh Wf T8 + \dots$$

$$rW4 = W4 T0 + \dots$$

$$rWf = \overline{W7} \overline{W8} W4 (T6 + T5) + \dots$$

$$Wc = Wh \overline{Wf} (T3 - T0) + \dots$$

3.168 IORP - $Iwh \overline{Iwi}$

3.169 Input

3.170 Read C words. If the channel counts C down to zero before the peripheral device encounters the End-of-Record (EOR), the channel ignores the rest of the record (to the End-of-Record). When the peripheral device sends the End-of-Record signal to the channel, the channel sets its End-of-Record Indicator; this signal sets the End-of-Record interrupt (if armed). The channel does not disconnect. The channel is now in an "Inter-Record" condition.

3. 171 When the peripheral device is magnetic tape, the tape continues to move when the tape handler encounters the End-of-Record. The End-of-Record occurs when the tape read-heads encounter tape gap; this also causes a Tape Gap signal to come high. If the program executes a new read tape or scan tape EOM during the inter-gap time (approximately one millisecond while the Tape Gap signal is high), the tape remains in motion and proceeds to read or scan the next record. If the program executes no such EOM before the Tape Gap signal drops, the channel disconnects and the tape comes to a stop. No additional interrupt occurs. This is the only condition that causes a channel to disconnect automatically in an IORP.

3. 172 All other input devices remain connected until the program takes further action. The paper tape reader remains in motion; the program should issue a "disconnect channel" instruction if the program is not reading any more tape. To proceed after the End-of-Record occurs, the program first executes a Buffer Control mode EOM to re-initialize the Channel Unit Address Register and then reloads the interlace portion of the channel (the program can alert the Interlace via the Buffer Control EOM). Otherwise, the channel immediately terminates any attempt to use its interlace portion since the channel is aware that it is still active and in the End-of-Record condition. When the program continues from an Inter-Record condition, the program should use an extended mode terminal function. An IORP should not be used to read with devices that do not have EOR signals (e.g., the typewriter).

3. 173 When the W-Register acquires the specified number of characters, a Time-Share Request, TRQW, is generated.

$$Trqw = \overline{Wf} \overline{W0} \overline{Wh} \overline{Iw} \overline{Iwf}$$

When memory is accessed, the interlace registers are counted by Interlace Count Trigger, Iwa.

$$Iwa = Rwx Tsm$$

Zero Word Count, Iwf, may occur and the Interlace Active flip-flop, Iw, is reset.

$$Iwf = Wc0 Wc1 \dots Wc13 Wc14 \overline{Ew}$$

$$rIw = Iwf T8 + \dots$$

If the Zero Word Count Interrupt Enable, Iwk, has been previously armed, then an Iiw interrupt occurs.

$$Iiw = Iwg \overline{Iw} Iwf Iwk + \dots$$

$$rIwk = Iiw Iwk T8 + \dots$$

3. 174 Additional characters entering the channel after Zero Word Count has been reached are precessed into the W-Register.

$$sW6 = \overline{W5} \overline{Ecw} T8 \overline{W10} \overline{W11} \overline{W12} \overline{W13} \overline{W14}$$

$$rW6 = W5 T0 + \dots$$

$$sW5 = \overline{W5} \overline{W6} \overline{Ecw} T0 + \dots$$

$$rW5 = W4 T0 + \dots$$

$$sW4 = W5 \overline{Wf} T8 \overline{Wg} + Iwg \overline{W9} \overline{Iwi} W5 \overline{Iwf} T8 + \dots$$

$$rW4 = W4 T0 + \dots$$

$$Ww1 = W4 Wb1 (T7 - T0) + \dots$$

$$Ww2 = W4 Wb2 (T7 - T0) + \dots$$

$$Ww3 = W4 Wb3 + \dots$$

However, Time Share Requests, Trqw, are inhibited.

$$Trqw = \overline{Wf} \overline{W0} \overline{Wh} \overline{Iw} \overline{Iwf}$$

Parity and rate errors cannot occur after Zero Word Count.

$$sWe = \overline{W9} \overline{W6} \overline{W5} \overline{W4} \overline{Rwp} \overline{Wg} \overline{Npw} (\overline{Iwg} + Iwi + \overline{Iwf}) + W0 \overline{W5} \overline{W6} \overline{Ecw} T8 + \dots$$

Detection of magnetic tape gap, photoreader gap or a halt signal sets the End-of-Record Detector.

$$sWg = Mtgw T0 Iwg W11 (W0 \overline{W9} + \dots) + Wbs (\overline{T7 - T0}) \overline{W10} \overline{W11} \overline{W12} \overline{W13} \overline{W14} + \overline{W9} \overline{W10} \overline{W11} \overline{W12} \overline{W13} (\overline{Rw1} \overline{Rw2} \overline{Rw3} \overline{Rw4} \overline{Rw5} \overline{Rw6} \overline{Rwp}) W5 (\overline{T7 - T0}) \overline{W10} \overline{W11} \overline{W12} \overline{W13} \overline{W14}$$

If the End-of-Record Detector is set before Zero Word Count has occurred, then the buffer is flushed and the completed word is stored in memory.

$$sW4 = Iwg Wg Wf W0 \overline{Wev} \overline{Iw} \overline{W7} \overline{W9} \overline{W10} \overline{W11} \overline{Wh} T8 + \dots$$

$$rW4 = W4 T0 + \dots$$

where:

$$Wev = W8 \overline{Wn2} \overline{W7} \overline{Wn1} + \overline{W8} \overline{Wn2} \overline{W7} \overline{Wn1} + \overline{W8} \overline{Wn2} \overline{W7} \overline{Wn1} + W8 \overline{Wn2} \overline{W7} \overline{Wn1}$$

Wf is reset and Time Share Request is inhibited.

$$rWf = \overline{W7} \overline{W8} W4 (T6 + T5) + \dots$$

$$Trqw = \overline{Wf} W0 \overline{Wh} Iw \overline{Iwf}$$

$$sWf = Iwa \overline{Mit} T0 + \dots$$

If Wg is set, Wsc is set after the buffer is flushed or immediately if Zero Word Count exists.

$$sWsc = Wg Iwg Iwh (Wew Wf + Iwf) T8 + \dots$$

$$rWsc = Wsc \overline{Iwg} T8$$

If the End-of-Record Interrupt Enable, Iwj, has been previously armed, an I2w interrupt occurs.

$$I2w = Wsc Iwj Iwg + \dots$$

$$rIwj = I2w Iwg T8 + \dots$$

3. 175 In general, the buffer does not disconnect, as the Halt Detector has not been set. For magnetic tape operation, a new EOM may be given within 0.75 millisecond from the occurrence of I2w to permit the magnetic tape system to proceed to a new record. In the case of magnetic tape, failure to give an EOM results in the tape stopping and the buffer disconnecting.

$$sWh = Whs W11 T8 + \dots$$

The buffer is flushed but not stored as Trqw is inhibited.

$$sW4 = Wh Wf T8 + \dots$$

$$rW4 = W4 T0 + \dots$$

$$rWf = \overline{W7} \overline{W8} W4 (T6 + T5)$$

$$Trqw = \overline{Wf} W0 \overline{Wh} Iw \overline{Iwf}$$

The buffer is cleared.

$$Wc = Wh \overline{Wf} (T3 - T0) + \dots$$

3. 176 IOSP - Iwg Iwh Iwi

3. 177 Input

3. 178 Read C words. If the channel counts C down to zero before the peripheral device encounters the End-of-Record, the channel generates a Count Equals Zero interrupt (if armed). The program should reload the interlace portion of the channel to continue reading the record. As far as the peripheral device knows, nothing happens at this time. Failure to reload the Interlace before the peripheral device sends enough characters to overflow the channel buffer causes a rate error; this sets the channel error indicator.

3. 179 When the peripheral device encounters the End-of-Record, IOSP operates identically like the IORP command. An IOSP is identical to an IORP in operation except that when Zero Word Count occurs it is anticipated that the interlace will be reloaded. Failure to reload the interlace in time results in rate errors. Parity error detection is not inhibited after Zero Word Count.

$$sWe = \overline{W9} \overline{W6} \overline{W5} \overline{W4} Rwp \overline{Wg} Npw (Iwi + \dots) + W0 \overline{W6} W5 Ecw T8 + \dots$$

3. 180 PIN ADDRESS COUNTER

3. 181 The PIN Address Counter flip-flop allows a PIN instruction to interrogate the Interlace Address Counter. The flip-flop is set by an I/O Unit Control instruction EOM (Ioc) in which the address bits are "zeros" and bit 13 is a "one". Bit 17 of the instruction selects either the W or Y channel TMCC.

$$sWpa = Ioc \overline{C17} C13 \overline{C19} \overline{C20} \overline{C21} \overline{C22} \overline{C23} T0$$

The Ioc signal indicates that an I/O control instruction is being processed and also distinguishes between the W - Y channels and the C - D channels.

$$Ioc = Ioc\ 1\ \overline{C1}\ \overline{Er}\ Qr3 \quad \overline{C1}\ \text{becomes } C1\ \text{for TMCC-C \& D}$$

$$Ioc\ 1 = Eom\ \overline{C10}\ C11$$

3. 182 After Wpa is turned on, the Interlace Memory Address Counter flip-flops are gated to the CPU via the parallel input lines (Cd0) - (Cd23). For this purpose, only fifteen of the lines are required.

$$(Cd9) = (Wa0\ Wpa)\ (Ya0\ Ypa) + \dots$$

⋮
⋮
⋮
⋮
⋮
⋮
⋮
⋮
⋮
⋮
⋮
⋮
⋮
⋮
⋮

$$(Cd23) = (Wa14\ Wpa)\ (Ya14\ Ypa) + \dots$$

The dotted lines indicate other parallel signals may be connected to the same lines.

3. 183 When a PIN instruction is executed, Wpa is re reset and the interrogation is complete.

$$rWpa = Pin\ T0 + St$$

3. 184 GLOSSARY OF LOGIC TERMS

		I2w	An interrupt signal indicating in the compatible mode that the input or output process has been terminated and the external device has been disconnected and in the extended mode indicating that an End-of-Record condition has been detected.
Buc	A control signal derived from the EOM instruction used to activate the TMCC and peripheral devices.		
C0 - C23	The 24 signals received from the C register in the CPU.	Ioc	An input/output control signal derived from Eom.
$\overline{C1x}$	A signal from the C1 flip-flop in the C register in the CPU. $\overline{C1x}$ is utilized to distinguish between TMCC W or Y and TMCC C or D. $\overline{C1x}$ becomes C1x for TMCC C or D.	$\overline{I1r0}/ - \overline{I1r14}/$	Interlace address signals transmitted to the CPU.
		Iw	A flip-flop which when set denotes that the interlace system is active.
C21r - C23r	Signals from the C register in the CPU used when serial transfers occur from the C register to the W register.	Iwa	A signal which counts the interlace word and address counters when a memory access is performed.
$\textcircled{Cd0} - \textcircled{Cd23}$	Input data lines that are read by the C register in the CPU during a PIN instruction.	Iwc	A signal occurring during the interlace loading sequence which clears the interlace registers.
$\overline{Ci0}/ - \overline{Ci23}/$	The 24 signals derived from the C register in the CPU and transmitted on the interconnecting cable	Iwe	A signal occurring during the EOM instruction of an interlace loading sequence which sets the high order bits of the interlace counters and set several flip-flops used in the extended mode.
Cpr	A signal from the CPU indicating the parity of C21r - C23r from the C register as information is serially transmitted to the TMCC.	Iwf	A signal denoting that the Interlace Word Count Register has reached zero word count.
Ecw	The clocking signal from external devices used in clocking of information into or out of the TMCC.	Iwg	A flip-flop which when set denotes extended mode operation.
En + \textcircled{En}	A signal from the CPU denoting that interrupt system is enabled.	Iwh, Iwi	Two flip-flops comprising the Channel Command Register used in the extended mode to control terminal functions of the TMCC.
Eom	A signal from the CPU occurring during an EOM instruction.	Iwj	The Channel Command Interrupt Enable flip-flop used in the extended mode to permit selective arming of the I2w interrupt signal.
Er	A signal which inhibits Eom and Ioc to external devices when an Interlace Prepare flip-flop has been set.	Iwk	The Channel Command Interrupt Enable flip-flop used in the extended mode to permit selective arming of the I1w interrupt signal.
Ew	A flip-flop which when set prepares the interlace to be loaded.		
I1w	An interrupt signal indicating in the compatible mode that a WIM or MIW instruction should either empty or load the TMCC and in the extended mode indicating that the interlace word count is zero.	Iwp	A signal occurring during the POT instruction of an interlace loading sequence which sets initial word count and memory address information into the Interlace Word and Address Counters.

$\overline{/Kcc0/ - /Kcc2/}$	Switch controlled signals from the Control Console used to select a particular TMCC or DACC for Unit Address and Error display on the Control Console.	$\overline{/Ra0/ - /Ra3/}$	Signals from the TMCC to the CPU denoting which particular TMCC's are active. These signals are displayed on the SDS 9300 Control Console and not used on the SDS 930.
$\overline{/Kcc1x/}$	A switch controlled signal used to distinguish between TMCC W or Y and TMCC C or D. $\overline{/Kcc1/}$ becomes $\overline{/Kcc1x/}$ for TMCC C or D.	$\overline{/Rd9/ - /Rd14/}$	Signals from the TMCC to the CPU denoting Unit Address Register contents for the particular TMCC selected for displays. These signals are displayed on the Control Console.
Kccw	A signal decoded from $\overline{/Kcc0/ - /Kcc2/}$ such that when true enables displaying the TMCC-W Unit Address and Error status.	$\overline{/Rde/}$	A signal from the TMCC to the CPU denoting the status of the Error Detector for the particular TMCC selected for display. This signal is displayed on the Control Console.
Mit	A signal derived from memory indicating the DACC is in process of accessing the memory.	$\overline{/Re0/ - /Re3/}$	Signals from the TMCC to the CPU denoting the Error Detector status for each particular TMCC. These signals are displayed on the SDS 9300 Control Console and not used on the SDS 930.
Mtgw	The magnetic tape gap signal generated by the magnetic tape system.		
Npw	A signal from external devices used to inhibit parity checking by the TMCC. When Npw is low, parity checking is inhibited.	$\overline{/Rr1/ - /Rr3/}$	Serial data signals transmitted from the TMCC to the CPU during time share operations.
05, Pwy	A signal from the CPU Operation Code Register, of the SDS 925 and 930. "05" is always true from the 9300.	\textcircled{Rt}	A ready signal from the TMCC on external devices used to release PIN or POT instructions from $\emptyset 2$.
Pin	A signal from the CPU derived from a PIN instruction.	Rti	A signal from the CPU to external devices indicating that a PIN instruction has terminated.
Pot	A signal from the CPU derived from a POT instruction.	Rw1-Rw24	Twenty four flip-flops comprising the character buffer. Rw1 through Rw6 are basic. Rw7 through Rw12 are optionally added to expand the character buffer to 12 bits. Rw7 through Rw24 are optionally added to expand the character buffer to 24 bits. TMCC Y, C, D are the only buffers which may be expanded.
Qr1-Qr4	The four flip-flops in the TMCC comprising the pulse counter.		
$\textcircled{Qq1}$	Timing signal sent to external devices which is true from T5 through T0.		
$\textcircled{Qq2}$	Timing signal sent to external devices which is true from T6 through T3.		
$\textcircled{Qq3}$	Timing signal sent to external devices which is true from T7 through T4.	Rwp	The parity flip-flop in the TMCC. An "odd" parity system is used.
Qw2	Timing signal which varies when a TMCC character length is expanded.	Rwx	A flip-flop in each TMCC which is set if a Time Share Request for the particular TMCC can occur.
$\overline{/R9/}$	A signal transmitted from the TMCC to the CPU defining whether the TMCC memory access operation is a load or store.	$\overline{/Rwy1/ - /Rwy3/}$	Serial data signals transmitted from the TMCC to the CPU during WIM instructions.

Rx	A signal from the CPU denoting that a WIM or MIW instruction is occurring.	W4	A flip-flop in the TMCC which controls the precessing of data between the Character Buffer and the W register.
Sio	A response signal from peripheral devices interrogated by an SKS instructions.	W5	A flip-flop which detects that a process should occur.
Skss	A signal generated in the CPU during SKS instructions and sent to external equipment to be used as a strobe.	W6	A flip-flop which detects that an external clock is present.
\sqrt{Skz}	A signal generated by TMCC's or DACC's which is sent to the CPU for interrogation during SKS instructions.	W7, W8	Two flip-flops comprising the Character Counter.
Ssc	A response signal from external systems equipment interrogated by SKS instructions.	W9	A flip-flop which is part of the Unit Address Register defining whether a process is input or output.
St	A signal from the CPU derived from the manual start button.	W10-W14	The Unit Address Register which designates to the TMCC and peripheral devices which device is activated.
Sys	A control signal for systems communication derived from EOM.	Wa0-Wa14	Fifteen flip-flops comprising the Interlace Address Register.
T8, T7-T0, T6 + T5, T6-T0, T4-T0, T3-T0, T0	Timing pulses used in the TMCC decoded from the Pulse Counter, Qr1 through Qr4.	Wb1-Wb3	Inputs to the W register from the Character Register which will vary depending on whether the TMCC is expanded.
Tpc	A timing pulse from the CPU used to synchronize the Pulse Counter, Qr1 through Qr4 in the TMCC with the Pulse Counter, Q1 through Q6, in the CPU.	Wc	The signal which resets the TMCC and prepares it for a new operation.
Trq	A signal transmitted from the TMCC to the CPU indicating that a TMCC requests a Time Share operation.	Wc0-Wc14	Fifteen flip-flops comprising the Interlace Word Counter.
Trqx	A signal transmitted from TMCC C and D indicating to TMCC W and Y that TMCC C and D are requesting a Time Share operation.	We	The Error Detector flip-flop.
Trqw	A signal generated by TMCC-W indicating that TMCC-W is requesting a Time Share operation.	Wes	An error signal from peripheral devices.
Tsm	A signal from the CPU indicating that the interlace address counter information has been received by the CPU.	Wev	A signal which when true indicates that the Character Counter is set to the same character count as it was when initially set up by the EOM.
Tsr	A signal from the CPU indicating that a Time Share operation is in process.	Wf	A flip-flop which when reset denotes on input that the W register is full and on output that the W register is empty.
		Wg	A flip-flop used to detect End-of-Record conditions.
		Wh	The Halt Detector flip-flop.

- Whs A halt signal from peripheral devices.
- Wn1-Wn3 The "now" flip-flops of the W register.
- W0 The Halt Interlock flip-flop used on output to enable the Halt Detector after output has been terminated and used on input to denote that the input process has proceeded to process characters.
- Wpa The PIN Address Counter flip-flop which when set allows a PIN instruction to interrogate the Interlace Address Counter.
- Wr1-Wr3 The "read" flip-flops of the W register.
- Ws The signal derived from an EOM instruction which initially sets up the TMCC.
- Wsc The Signal Complete flip-flop.
- Ww1-Ww3 The signals which permit "writing" into the W register.
- Wxx A signal which denotes that an MIW, WIM or Time Share operation is occurring.
- Zw1-Zw24 The twenty four inputs to the Character Register from peripheral devices.
- Zwp The input to the Parity flip-flop from peripheral devices.

$$\begin{aligned}
 T8 &= Qr1 \overline{Qr3} \overline{Qr4} \\
 T6 + T5 &= Qr2 \overline{Qr4} \\
 T7 - T4 &= Qr4 \\
 T7 - T0 &= Qr3 + Qr4 \\
 T6 - T0 &= Qr3 \\
 T6 - T3 &= (Qr4 + \overline{Qr1} \overline{Qr2}) Qr3 \\
 T3 - T0 &= Qr3 \overline{Qr4} \\
 T0 &= \overline{Qr1} \overline{Qr2} \overline{Qr4} \\
 T5 - T1 &= Qr3 \overline{Qr1} \overline{Qr4} \overline{T0} \\
 T5 - T0 &= Qr3 \overline{Qr1} \overline{Qr4}
 \end{aligned}$$

3. 185 LOGIC EQUATIONS

3. 186 Pulse Counter

$$sQr1 = Tpc + \overline{Qr2} \overline{Qr3} \overline{Qr4}$$

$$rQr1 = Qr2$$

$$sQr2 = Qr1 \overline{Qr2} \overline{Qr4} + (Qr4 + T0) + Qr1 \overline{Qr3}$$

$$rQr2 = \overline{Qr1}$$

$$sQr3 = Qr1 \overline{Qr4}$$

$$rQr3 = Qr3 \overline{Qr4} (Qr4 + T0)$$

$$sQr4 = Qr1 \overline{Qr3}$$

$$rQr4 = \overline{Qr1} \overline{Qr2}$$

	Qr1	Qr2	Qr3	Qr4
Tp	0	0	0	0
T8	1	0	0	0
T7	1	0	0	1
T6	1	1	1	1
T5	0	1	1	1
T4	0	0	1	1
T3	0	0	1	0
T2	1	0	1	0
T1	1	1	1	0
T0	0	1	1	0
Tr	0	0	0	0
Tp	0	0	0	0

3. 187 Buc, Ioc, Sys, etc.

$$Buc = Eom \overline{C10} \overline{C11} \overline{C1}^*$$

$$Ioc1 = Eom \overline{C10} C11$$

$$Ioc = Ioc1 \overline{C1} \overline{Er} \overline{Qr3}^*$$

$$Sys = Eom C10 C11 \overline{C9}$$

* $\overline{C1}$ becomes C1 for TMCC-C and TMCC-D

3. 188 CPU Signals Received

C0	=	$\overline{Ci0}$
C1	=	
C2	=	
C3	=	
C4	=	
C5	=	
C6	=	
C7	=	
C8	=	
C9	=	
C10	=	
C11	=	
C12	=	
C13	=	
C14	=	
C15	=	
C16	=	
C17	=	

C18	=	$\sqrt{Ci18}$
C19	=	-----
C20	=	-----
C21	=	-----
C22	=	-----
C23	=	$\sqrt{Ci23}$
\overline{Tpc}	=	\sqrt{Tpc}
Eom	=	\sqrt{Eom}
Pot 1	=	\sqrt{Pot}
Pin	=	\sqrt{Pin}
\overline{Rti}	=	\sqrt{Rti}
Cpr	=	\sqrt{Cpr}
C21r	=	$\sqrt{C21r}$
C22r	=	$\sqrt{C22r}$
C23r	=	$\sqrt{C23r}$
Rx	=	\sqrt{Rx}
\overline{Tsm}	=	\sqrt{Tsm}
Tsr	=	\sqrt{Tsr}
C1x	=	$\sqrt{C1x}$ *For TMCC-C and TMCC-D: C1x = $\sqrt{C1x}$
Pwy	=	\sqrt{Pwy} (Pwy = 05 for 92200 and 92210)
\overline{Mit}	=	\sqrt{Mit}
En + \textcircled{En}	=	$\sqrt{En + \textcircled{En}}$ /
St	=	\sqrt{St}
Skss	=	\sqrt{Skss}
Kccw	=	$\sqrt{Kcc0/ + /Kcc1/ + /Kcc2/*}$
$\overline{Kcc2}$	=	$\sqrt{Kcc2/}$
Kccy	=	$\sqrt{Kcc0/ + /Kcc1/ + /Kcc2/*}$

*For TMCC-C and TMCC-D, $\sqrt{Kcc1/}$ becomes $\sqrt{\overline{Kcc1/}}$

3.189 Input/Output Signals Generated

$\textcircled{C0}$	=	C0
$\textcircled{C1}$	=	C1
$\textcircled{C2}$	=	C2
$\textcircled{C3}$	=	C3
$\textcircled{C4}$	=	C4
$\textcircled{C5}$	=	C5
$\textcircled{C6}$	=	C6

$\textcircled{C7}$	=	C7
$\textcircled{C8}$	=	C8
$\textcircled{C9}$	=	C9
$\textcircled{C10}$	=	C10
$\textcircled{C11}$	=	C11
$\textcircled{C12}$	=	C12
$\textcircled{C13}$	=	C13
$\textcircled{C14}$	=	C14
$\textcircled{C15}$	=	C15
$\textcircled{C16}$	=	C16
$\textcircled{C17}$	=	C17
$\textcircled{C18}$	=	C18
$\textcircled{C19}$	=	C19
$\textcircled{C20}$	=	C20
$\textcircled{C21}$	=	C21
$\textcircled{C22}$	=	C22
$\textcircled{C23}$	=	C23
$\textcircled{C17}$	=	C17
$\textcircled{C24}$	=	Cpr (93200 and 92331 only)
$\textcircled{Qq1}$	=	Qr3 ($\overline{Qr1}$ $\overline{Qr4}$) = T5 - T0
$\textcircled{Qq2}$	=	(T7 - T3) Qr3 = T6 - T3
$\textcircled{Qq3}$	=	Qr4 = T7 - T4 Replaces Mtgw on POT connectors at 23F & 24F
\textcircled{Buc}	=	Buc
\textcircled{Ioc}	=	Ioc
\textcircled{Eom}	=	\overline{Eom} \overline{Er}
\textcircled{Sys}	=	Sys (T5 - T1) \overline{Tsr} Pwy (delete Pwy for 92200 and 92210)
\textcircled{St}	=	St
\textcircled{Pin}	=	Pin
$\textcircled{Pot 1}$	=	Pot 1
$\textcircled{Pot 2}$	=	Pot 1 (T5 - T1) \overline{Tsr} Pwy (delete Pwy for 92200 and 92210)
\textcircled{Rti}	=	\overline{Rti}
\textcircled{Skss}	=	Skss

3.190 CPU Signals Generated

$\sqrt{Rd9/}$	=	$\overline{W9}$ \overline{Kccw} $\overline{Y9}$ \overline{Kccy} . . .
$\sqrt{Rd10/}$	=	$\overline{W10}$ \overline{Kccw} $\overline{Y10}$ \overline{Kccy} . . .

$\overline{Rd11}/$	=	$\overline{W11 Kccw Y11 Kccy} \dots$
$\overline{Rd12}/$	=	$\overline{W12 Kccw Y12 Kccy} \dots$
$\overline{Rd13}/$	=	$\overline{W13 Kccw Y13 Kccy} \dots$
$\overline{Rd14}/$	=	$\overline{W14 Kccw Y14 Kccy} \dots$
$\overline{Rde}/$	=	$\overline{We Kccw Ye Kccy} \dots$
$\overline{Ra0}/$	=	$\overline{W10 W11 W12 W13 W14}$ $\overline{Ra2}/$ for TMCC-C
$\overline{Ra1}/$	=	$\overline{Y10 Y11 Y12 Y13 Y14}$ $\overline{Ra3}/$ for TMCC-D
$\overline{Re0}/$	=	We $\overline{Re2}/$ for TMCC-C
$\overline{Re1}/$	=	Ye $\overline{Re3}/$ for TMCC-D
$\overline{Wf (W0+W9)}/$	=	$\overline{Wf (W0+W9)}$
$\overline{Yf (Y0+Y9)}/$	=	$\overline{Yf (Y0+Y9)}$
$\overline{Trq}/$	=	$\overline{Trqw Trqy Trqx}$
$\overline{I1w}/$	=	$\overline{I1w}$ $\overline{I1c}/$ for TMCC-C
$\overline{I2w}/$	=	$\overline{I2w}$ $\overline{I2c}/$ for TMCC-C
$\overline{I1y}/$	=	$\overline{I1y}$ $\overline{I1d}/$ for TMCC-D
$\overline{I2y}/$	=	$\overline{I2y}$ $\overline{I2d}/$ for TMCC-D
$\overline{Skrz}/$	=	\overline{Skr} $\overline{C10 C11 Ssc C1 C9}$ $\overline{C10 C11 Sio} \dots *$

* $\overline{C1}$ becomes C1 for TMCC-C and TMCC-D

$\overline{Rr1}/$	=	$\overline{Wr1 Rwx Yr1 Ryx} \dots$
$\overline{Rr2}/$	=	$\overline{Wr2 Rwx Yr2 Ryx} \dots$
$\overline{Rr3}/$	=	$\overline{Wr3 Rwx Yr3 Ryx} \dots$
$\overline{Rwy1}/$	=	$\overline{Wr1 Pwy Yr1 Pwy}$ (Pwy = 05 for 92200 and 92210)
$\overline{Rwy2}/$	=	$\overline{Wr2 Pwy Yr2 Pwy}$ (Pwy = 05 for 92200 and 92210)
$\overline{Rwy3}/$	=	$\overline{Wr3 Pwy Yr3 Pwy}$ (Pwy = 05 for 92200 and 92210)

$\overline{Cd0}$	=	$\overline{Cd0} \dots$
$\overline{Cd1}$	=	$\overline{Cd1} \dots$
$\overline{Cd2}$	=	$\overline{Cd2} \dots$
$\overline{Cd3}$	=	$\overline{Cd3} \dots$
$\overline{Cd4}$	=	$\overline{Cd4} \dots$
$\overline{Cd5}$	=	$\overline{Cd5} \dots$
$\overline{Cd6}$	=	$\overline{Cd6} \dots$

$\overline{Cd7}$	=	$\overline{Cd7} \dots$
$\overline{Cd8}$	=	$\overline{Cd8} \dots$
$\overline{Cd9}$	=	$\overline{Cd9 Wa0 Wpa Ya0 Ypa} \dots$
$\overline{Cd10}$	=	$\overline{Cd10 Wa1 Wpa Ya1 Ypa} \dots$
$\overline{Cd11}$	=	$\overline{Cd11 Wa2 Wpa Ya2 Ypa} \dots$
$\overline{Cd12}$	=	$\overline{Cd12 Wa3 Wpa Ya3 Ypa} \dots$
$\overline{Cd13}$	=	$\overline{Cd13 Wa4 Wpa Ya4 Ypa} \dots$
$\overline{Cd14}$	=	$\overline{Cd14 Wa5 Wpa Ya5 Ypa} \dots$
$\overline{Cd15}$	=	$\overline{Cd15 Wa6 Wpa Ya6 Ypa} \dots$
$\overline{Cd16}$	=	$\overline{Cd16 Wa7 Wpa Ya7 Ypa} \dots$
$\overline{Cd17}$	=	$\overline{Cd17 Wa8 Wpa Ya8 Ypa} \dots$
$\overline{Cd18}$	=	$\overline{Cd18 Wa9 Wpa Ya9 Ypa} \dots$
$\overline{Cd19}$	=	$\overline{Cd19 Wa10 Wpa Ya10 Ypa} \dots$
$\overline{Cd20}$	=	$\overline{Cd20 Wa11 Wpa Ya11 Ypa} \dots$
$\overline{Cd21}$	=	$\overline{Cd21 Wa12 Wpa Ya12 Ypa} \dots$
$\overline{Cd22}$	=	$\overline{Cd22 Wa13 Wpa Ya13 Ypa} \dots$
$\overline{Cd23}$	=	$\overline{Cd23 Wa14 Wpa Ya14 Ypa} \dots$
\overline{Rf}	=	$\overline{Rf Wap + Ew Yap + Ey} \dots$
$\overline{Cd24}$	=	$\overline{Cd24}$ (93200 and 93221 only)

\overline{Bf} (93200 and 93221 only)

$\overline{Ir0}/$	=	$\overline{Wa0 Rwx Ya0 Ryx} \dots$
$\overline{Ir1}/$	=	$\overline{Wa1 Rwx Ya1 Ryx} \dots$
$\overline{Ir2}/$	=	$\overline{Wa2 Rwx Ya2 Ryx} \dots$
$\overline{Ir3}/$	=	$\overline{Wa3 Rwx Ya3 Ryx} \dots$
$\overline{Ir4}/$	=	$\overline{Wa4 Rwx Ya4 Ryx} \dots$
$\overline{Ir5}/$	=	$\overline{Wa5 Rwx Ya5 Ryx} \dots$
$\overline{Ir6}/$	=	$\overline{Wa6 Rwx Ya6 Ryx} \dots$
$\overline{Ir7}/$	=	$\overline{Wa7 Rwx Ya7 Ryx} \dots$
$\overline{Ir8}/$	=	$\overline{Wa8 Rwx Ya8 Ryx} \dots$
$\overline{Ir9}/$	=	$\overline{Wa9 Rwx Ya9 Ryx} \dots$
$\overline{Ir10}/$	=	$\overline{Wa10 Rwx Ya10 Ryx} \dots$
$\overline{Ir11}/$	=	$\overline{Wa11 Rwx Ya11 Ryx} \dots$
$\overline{Ir12}/$	=	$\overline{Wa12 Rwx Ya12 Ryx} \dots$
$\overline{Ir13}/$	=	$\overline{Wa13 Rwx Ya13 Ryx} \dots$
$\overline{Ir14}/$	=	$\overline{Wa14 Rwx Ya14 Ryx} \dots$
$\overline{R9}/$	=	$\overline{W9 Rwx Y9 Ryx} \dots$

3. 191 TMCC Signals Received

$$Er = \sqrt{Er}/$$

$$Trqx = \sqrt{Trqx}/$$

3. 192 Input/Output Signals Received

$$Sio = \overline{Sio}$$

$$Ssc = \overline{Ssc}$$

3. 193 TMCC Signals Generated

$$\sqrt{Er}/ = \overline{Ew} \overline{Ey} . . .$$

$$\sqrt{Trqx}/ = \overline{Trq(c)} \overline{Trq(d)}$$

3. 194 LOGIC EQUATIONS FOR W BUFFER

3. 195 Unit Address Register

$$sW14 = Ws C23$$

$$rW14 = Wc$$

$$sW13 = Ws C22$$

$$rW13 = Wc$$

$$sW12 = Ws C21$$

$$rW12 = Wc$$

$$sW11 = Ws C20$$

$$rW11 = Wc$$

$$sW10 = Ws C19 + (Ioc C12 \overline{C17} \overline{C19} \overline{C20} \overline{C21} \overline{C22} \overline{C23}) \overline{W9} \overline{W10}$$

$$rW10 = Wc$$

3. 196 Input/Output

$$sW9 = Ws C18$$

$$rW9 = Wc$$

3. 197 Clear and Set Signals

$$Wc = Buc \overline{C17} (T6 + T5) + Wh \overline{Wf} (T3 - T0) + St$$

$$Ws = Buc \overline{C17} (T3 - T0)$$

3. 198 Clock Counter

$$sW6 = \overline{W5} Ecw T8 \overline{W10} \overline{W11} \overline{W12} \overline{W13} \overline{W14}$$

$$rW6 = W5 T0 + Wc$$

$$sW5 = \overline{W5} W6 \overline{Ecw} T0 + Ws C13 C18 T0 \overline{W10} \overline{W11} \overline{W12} \overline{W13} \overline{W14}$$

$$rW5 = W4 T0 + Wc$$

$$sW4 = W5 Wf T8 \overline{Wg} + Ws T0 + Wh Wf T8 + Iwg \overline{W9} \overline{Iwi} W5 Iwf T8 + Iwg Wg Wf W0 \overline{Wev} Iw T8 \overline{W7} \overline{W9} \overline{W10} \overline{W11} \overline{Wh}$$

$$rW4 = W4 T0 + W4 T8$$

3. 199 Character Counter

$$sW8 = Ws C16 + W7 \overline{W8} W4 T0 + Wxx Wn2 (\overline{T7} - T0) \overline{W4} + \overline{W7} \overline{W9} W10 W11 \overline{Wh}$$

$$rW8 = W8 W4 T0 + Wc$$

$$sW7 = Ws C15 + Wxx Wn1 (\overline{T7} - T0) \overline{W4}$$

$$rW7 = Wc + W7 \overline{W8} W4 T0$$

3. 200 Character Counter Even

$$Wev = W8 Wn2 W7 Wn1 + \overline{W8} \overline{Wn2} W7 Wn1 + \overline{W8} \overline{Wn2} \overline{W7} \overline{Wn1} + W8 Wn2 \overline{W7} \overline{Wn1}$$

3. 201 Halt Interlock

$$sW0 = \overline{W9} W6 \overline{W8} Ecw + Ws C18 W9$$

$$rW0 = Wc + (Ioc C12 \overline{C17} \overline{C19} \overline{C20} \overline{C21} \overline{C22} \overline{C23}) W9 T0 + W9 Iw Iwf (\overline{Iwg} + \overline{Iwh} + \overline{Iwi}) \overline{T7} - T0$$

3. 202 Computer Interlock

$$sWf = Wc \overline{Wh} + Iwa \overline{Mit} T0 + Rx T0 Pwy *(Pwy = 05 for 92200)$$

*Rx is always false for TMCC-C

$$\begin{aligned}
 rWf &= \overline{W7} \overline{W8} W4 (T6 + T5) \\
 &+ Ws C18 \\
 &+ \overline{W9} W10 W11 W0 Mtgw \overline{W7} (T6 + T5) \overline{Wh}
 \end{aligned}$$

3.203 End-of-Record Detector

$$\begin{aligned}
 sWg &= Mtgw T0 Iwg W11 (\overline{W0} \overline{W9} + \overline{W0} W5 \overline{W6} W9) \\
 &+ Whs \overline{T7} - \overline{T0} \overline{W10} \overline{W11} \overline{W12} \overline{W13} \overline{W14} \\
 &+ \overline{W9} \overline{W10} \overline{W11} W12 \overline{W13} (\overline{Rw1} \overline{Rw2} \overline{Rw3} \\
 &\quad \overline{Rw4} \overline{Rw5} \overline{Rw6} \overline{Rwp}) W5 \overline{T7} - \overline{T0} \\
 &\quad \overline{W10} \overline{W11} \overline{W12} \overline{W13} \overline{W14} \\
 rWg &= Wc \\
 &+ Wg \overline{Iwg} W11 T0
 \end{aligned}$$

3.204 Halt Detector

$$\begin{aligned}
 sWh &= Wg \overline{Iwg} T8 \\
 &+ Whs W11 T8 \\
 &+ Wg \overline{W9} \overline{W11} \overline{Iwh} (Iwf + Wev Wf) T8 Iwg \\
 &+ W9 \overline{W11} \overline{W0} W5 \overline{W6} (\overline{Iwg} + Iwi \overline{Iwh}) T8 \\
 &+ \overline{W9} \overline{W11} Iwg \overline{Iwh} Iwi Iwf T8 W0 \\
 &+ W9 \overline{Iwh} T8 Whs \\
 rWh &= Wc (T6 + T5) \\
 &+ Wh \overline{Wf} T8
 \end{aligned}$$

3.205 Signal Complete

$$\begin{aligned}
 sWsc &= Wg Iwg Iwh (Wev Wf + Iwf) T8 \\
 &+ Wh \overline{Wf} T8 + St \\
 &+ Ws \overline{C19} \overline{C20} \overline{C21} \overline{C22} \overline{C23} \\
 rWsc &= Wsc T8 \overline{Iwg}
 \end{aligned}$$

3.206 Interrupt Signals

$$\begin{aligned}
 I1w &= \overline{Wf} W0 \overline{Wh} (En + \textcircled{En}) \overline{Iw} \overline{Ew} \overline{Iwg} \\
 &+ Iwg \overline{Iw} Iwf Iwk \\
 I2w &= (En + \textcircled{En}) \overline{Iwg} Wh \overline{Wf} \\
 &+ Wsc Iwj Iwg
 \end{aligned}$$

3.207 WIM + MIW Interlock =

$$\overline{Wf} (W0 + \overline{W9})$$

3.208 Load Buffer from C

$$\begin{aligned}
 Wxx &= Rx Pwy + Rwx Tsr Iw* \\
 &\quad (Pwy = 05 \text{ and } Iw \text{ deleted from } 92200)
 \end{aligned}$$

*Rx is always false for TMCC-C

3.209 Time Share Request

$$Trqw = \overline{Wf} W0 \overline{Wh} Iw \overline{Iwf}$$

3.210 Time Share Select

$$\begin{aligned}
 sRwx &= \overline{Tsm} Trqw \overline{Trqy} \overline{Trqx} (\overline{T7} - \overline{T0}) \\
 rRwx &= \overline{Tsm} T0
 \end{aligned}$$

3.211 Time Share Priority

$$\begin{aligned}
 Trqx &= \text{Priority signal from TMCC-C and TMCC-D} \\
 Trqx &= Trq(c) + Trq(d) \\
 Trq &= Trqx + Trqw + Trqy
 \end{aligned}$$

3.212 W Register

$$\begin{aligned}
 Ww1 &= W4 W7 (\overline{T7} - \overline{T0}) \\
 &+ \overline{W4} Wn1 (\overline{T7} - \overline{T0}) \\
 &+ \overline{W4} Wn1 \overline{Wxx} \\
 &+ W4 Wb1 (\overline{T7} - \overline{T0}) * \\
 &+ \overline{W4} C21r (\overline{T7} - \overline{T0}) Wxx \\
 Ww2 &= W4 W8 (\overline{T7} - \overline{T0}) \\
 &+ \overline{W4} Wn2 (\overline{T7} - \overline{T0}) \\
 &+ \overline{W4} Wn2 \overline{Wxx} \\
 &+ W4 Wb2 (\overline{T7} - \overline{T0}) * \\
 &+ \overline{W4} C22r (\overline{T7} - \overline{T0}) Wxx
 \end{aligned}$$

*sWr1, 2, 3 = Ww1, 2, 3 Delayed by 9 pulse times

sWn1, 2, 3 = Wr1, 2, 3 respectively

$$\left. \begin{aligned}
 * Wb1 &= \overline{Wx12} \overline{Wx24} Rw4 + Wx12 Rw10 \\
 &\quad + Wx24 Rw22 \\
 Wb2 &= \overline{Wx12} \overline{Wx24} Rw5 + Wx12 Rw11 \\
 &\quad + Wx24 Rw23
 \end{aligned} \right\} \text{for 93200 and 93221}$$

$$\begin{aligned} Ww3 &= \overline{W4} Wn3 \overline{Wxx} \\ &+ W4 Wb3 * \\ &+ \overline{W4} C23r Wxx \end{aligned}$$

* $Wb3 = \overline{Wx12} \overline{Wx24} Rw6 + Wx12 Rw12$ for 93200
 + $Wx24 Rw24$ and 93221

Wb1, 2, 3 equals Rw4, 5, 6 respectively for 92200
 Wb1, 2, 3 equals Rw10, 11, 12 respectively for 92201
 Wb1, 2, 3 equals Rw22, 23, 24 respectively for 92202

3. 213 Character Buffer*

$$\begin{aligned} sRw1 &= W4 \overline{Wxx} (T7 - T0) Wn1 + \overline{W9} W6 \overline{W5} Zw1 + \\ &W4 Wxx C21r \\ rRw1 &= W4 \overline{Wxx} \quad \overline{Wn1} + \overline{W9} \overline{W6} \overline{W5} \overline{W4} + \\ &W4 Wxx \overline{C21r} \end{aligned}$$

$$\begin{aligned} sRw2 &= W4 \overline{Wxx} (T7 - T0) Wn2 + \overline{W9} W6 \overline{W5} Zw2 + \\ &W4 Wxx C22r \\ rRw2 &= W4 \overline{Wxx} \quad \overline{Wn2} + \overline{W9} \overline{W6} \overline{W5} \overline{W4} + \\ &W4 Wxx \overline{C22r} \end{aligned}$$

$$\begin{aligned} sRw3 &= W4 \overline{Wxx} (T7 - T0) Wn3 + \overline{W9} W6 \overline{W5} \\ &Zw3 + W4 Wxx C23r \\ rRw3 &= W4 \overline{Wxx} \quad \overline{Wn3} + \overline{W9} \overline{W6} \overline{W5} \overline{W4} + \\ &W4 Wxx \overline{C23r} \end{aligned}$$

$$\begin{aligned} sRw4 &= W4 \overline{Rw4} Rw1 + \overline{W9} W6 \overline{W5} Zw4 \\ rRw4 &= W4 Rw4 \overline{Rw1} + \overline{W9} \overline{W6} \overline{W5} \overline{W4} \end{aligned}$$

$$\begin{aligned} sRw5 &= W4 \overline{Rw5} Rw2 + \overline{W9} W6 \overline{W5} Zw5 \\ rRw5 &= W4 Rw5 \overline{Rw2} + \overline{W9} \overline{W6} \overline{W5} \overline{W4} \end{aligned}$$

$$\begin{aligned} sRw6 &= W4 \overline{Rw6} Rw3 + \overline{W9} W6 \overline{W5} Zw6 \\ rRw6 &= W4 Rw6 \overline{Rw3} + \overline{W9} \overline{W6} \overline{W5} \overline{W4} \end{aligned}$$

*For 12-bit extension, add Rw7 through Rw12
 For 24-bit extension, add Rw7 through Rw24

Character Size

12 bit character = $Wx12$ (93200 and 93221 only)
 24 bit character = $Wx24$ (93200 and 93221 only)

3. 214 Character Buffer Extended to 12 Bits

$$\begin{aligned} sRw7 &= W4 \overline{Rw7} Rw4 + \overline{W9} W6 \overline{W5} Zw7 \\ rRw7 &= W4 Rw7 \overline{Rw4} + \overline{W9} \overline{W6} \overline{W5} \overline{W4} \end{aligned}$$

$$\begin{aligned} sRw8 &= W4 \overline{Rw8} Rw5 + \overline{W9} W6 \overline{W5} Zw8 \\ rRw8 &= W4 Rw8 \overline{Rw5} + \overline{W9} \overline{W6} \overline{W5} \overline{W4} \end{aligned}$$

$$\begin{aligned} sRw9 &= W4 \overline{Rw9} Rw6 + \overline{W9} W6 \overline{W5} Zw9 \\ rRw9 &= W4 Rw9 \overline{Rw6} + \overline{W9} \overline{W6} \overline{W5} \overline{W4} \end{aligned}$$

$$\begin{aligned} sRw10 &= W4 \overline{Rw10} Rw7 + \overline{W9} W6 \overline{W5} Zw10 \\ rRw10 &= W4 Rw10 \overline{Rw7} + \overline{W9} \overline{W6} \overline{W5} \overline{W4} \end{aligned}$$

$$\begin{aligned} sRw11 &= W4 \overline{Rw11} Rw8 + \overline{W9} W6 \overline{W5} Zw11 \\ rRw11 &= W4 Rw11 \overline{Rw8} + \overline{W9} \overline{W6} \overline{W5} \overline{W4} \end{aligned}$$

$$\begin{aligned} sRw12 &= W4 \overline{Rw12} Rw9 + \overline{W9} W6 \overline{W5} Zw12 \\ rRw12 &= W4 Rw12 \overline{Rw9} + \overline{W9} \overline{W6} \overline{W5} \overline{W4} \end{aligned}$$

3. 215 Character Buffer Extended to 24 Bits

$$\begin{aligned} sRw13 &= W4 \overline{Rw13} Rw10 + \overline{W9} W6 \overline{W5} Zw13 \\ rRw13 &= W4 Rw13 \overline{Rw10} + \overline{W9} \overline{W6} \overline{W5} \overline{W4} \end{aligned}$$

$$\begin{aligned} sRw14 &= W4 \overline{Rw14} Rw11 + \overline{W9} W6 \overline{W5} Zw14 \\ rRw14 &= W4 Rw14 \overline{Rw11} + \overline{W9} \overline{W6} \overline{W5} \overline{W4} \end{aligned}$$

$$\begin{aligned} sRw15 &= W4 \overline{Rw15} Rw12 + \overline{W9} W6 \overline{W5} Zw15 \\ rRw15 &= W4 Rw15 \overline{Rw12} + \overline{W9} \overline{W6} \overline{W5} \overline{W4} \end{aligned}$$

$$\begin{aligned} sRw16 &= W4 \overline{Rw16} Rw13 + \overline{W9} W6 \overline{W5} Zw16 \\ rRw16 &= W4 Rw16 \overline{Rw13} + \overline{W9} \overline{W6} \overline{W5} \overline{W4} \end{aligned}$$

$$\begin{aligned} sRw17 &= W4 \overline{Rw17} Rw14 + \overline{W9} W6 \overline{W5} Zw17 \\ rRw17 &= W4 Rw17 \overline{Rw14} + \overline{W9} \overline{W6} \overline{W5} \overline{W4} \end{aligned}$$

$$\begin{aligned} sRw18 &= W4 \overline{Rw18} Rw15 + \overline{W9} W6 \overline{W5} Zw18 \\ rRw18 &= W4 Rw18 \overline{Rw15} + \overline{W9} \overline{W6} \overline{W5} \overline{W4} \end{aligned}$$

$$\begin{aligned} sRw19 &= W4 \overline{Rw19} Rw16 + \overline{W9} W6 \overline{W5} Zw19 \\ rRw19 &= W4 Rw19 \overline{Rw16} + \overline{W9} \overline{W6} \overline{W5} \overline{W4} \end{aligned}$$

$$\begin{aligned} sRw20 &= W4 \overline{Rw20} Rw17 + \overline{W9} W6 \overline{W5} Zw20 \\ rRw20 &= W4 Rw20 \overline{Rw17} + \overline{W9} \overline{W6} \overline{W5} \overline{W4} \end{aligned}$$

$$\begin{aligned} sRw21 &= W4 \overline{Rw21} Rw18 + \overline{W9} W6 \overline{W5} Zw21 \\ rRw21 &= W4 Rw21 \overline{Rw18} + \overline{W9} \overline{W6} \overline{W5} \overline{W4} \end{aligned}$$

$$\begin{aligned} sRw22 &= W4 \overline{Rw22} Rw19 + \overline{W9} W6 \overline{W5} Zw22 \\ rRw22 &= W4 Rw22 \overline{Rw19} + \overline{W9} \overline{W6} \overline{W5} \overline{W4} \end{aligned}$$

$$sRw23 = W4 \overline{Rw23} \overline{Rw20} + \overline{W9} \overline{W6} \overline{W5} \overline{Zw23}$$

$$rRw23 = W4 \overline{Rw23} \overline{Rw20} + \overline{W9} \overline{W6} \overline{W5} \overline{W4}$$

$$sRw24 = W4 \overline{Rw24} \overline{Rw21} + \overline{W9} \overline{W6} \overline{W5} \overline{Zw24}$$

$$rRw24 = W4 \overline{Rw24} \overline{Rw21} + \overline{W9} \overline{W6} \overline{W5} \overline{W4}$$

3.216 Parity Flip-Flop

$$\begin{aligned} sRwp &= \overline{W9} \overline{W4} \overline{Rwp} (Wb1 \oplus Wb2 \oplus Wb3) \\ &\quad (T7 - T0) \overline{Qw1} * \\ &+ \overline{W9} \overline{W4} \overline{Rwp} \overline{Wxx} (Wn1 \oplus Wn2 \oplus Wn3) \\ &\quad \overline{Qw2} (T7 - T0) * \\ &+ \overline{W9} \overline{W4} \overline{Rwp} \overline{Wxx} \overline{Cpr} \overline{Qw2} (T7 - T0) * \\ &+ \overline{W9} \overline{W6} \overline{W5} \overline{Zwp} \\ &+ \overline{Wf} \overline{W5} \overline{T8} \overline{Rwp} \end{aligned}$$

$$\begin{aligned} rRwp &= \overline{W9} \overline{W4} \overline{Rwp} (Wb1 \oplus Wb2 \oplus Wb3) \\ &\quad (T7 - T0) \overline{Qw1} * \\ &+ \overline{W9} \overline{W4} \overline{Rwp} \overline{Wxx} (Wn1 \oplus Wn2 \oplus Wn3) \\ &\quad \overline{Qw2} (T7 - T0) * \\ &+ \overline{W9} \overline{W4} \overline{Rwp} \overline{Wxx} \overline{Cpr} \overline{Qw2} (T7 - T0) * \\ &+ \overline{W9} \overline{W6} \overline{W5} \overline{W4} \\ &+ \overline{Wf} \overline{W5} \overline{T8} \overline{W9} \overline{Rwp} \\ &+ \overline{Wc} \end{aligned}$$

$$* \overline{Qw1} = \overline{Wx12} \overline{Qr4} + \overline{Wx12} \overline{Wx24} \overline{Qr1} \overline{Qr4}$$

(for 93200 and 93221)

$$\overline{Qw2} = \overline{Wx12} \overline{Qr4} + \overline{Wx12} \overline{Wx24} \overline{Qr2} \overline{Qr4}$$

(for 93200 and 93221)

$$\overline{Qw1} = \overline{Qr1} \overline{Qr4} \text{ for } 92200$$

$$\overline{Qw1} = \overline{Qr4} \text{ for } 92201$$

$\overline{Qw1}$ is deleted for 92202

$$\overline{Qw2} = \overline{Qr2} \overline{Qr4} \text{ for } 92200$$

$$\overline{Qw2} = \overline{Qr4} \text{ for } 92201$$

$\overline{Qw2}$ is deleted for 92202

3.217 Error Detector

$$\begin{aligned} sWe &= \overline{W9} \overline{W6} \overline{W5} \overline{W4} \overline{Rwp} \overline{Wg} \overline{Npw} \\ &\quad (\overline{Iwg} + \overline{Iwi} + \overline{Iwf}) \\ &+ \overline{W0} \overline{W6} \overline{W5} \overline{Ecw} \overline{T8} \\ &+ \overline{Wes} \end{aligned}$$

$$rWe = \overline{Wc} \overline{Wh}$$

3.218 Interlace Prepare

$$sEw = \overline{Iwc} \overline{Ew} (T3 - T0)$$

$$rEw = \overline{Wc} \overline{T0}$$

$$+ \text{Pot } 1 (T3 - T0) \overline{Ew}$$

3.219 Interlace Clear

$$\overline{Iwc} = \overline{Eom} \overline{C9} \overline{C10} \overline{C1} \overline{C17} (T3 - T0) *$$

* $\overline{C1}$ becomes $C1$ for TMCC-C

3.220 Interlace Load

$$\overline{Iwp} = \text{Pot } 1 (T6 - T5) \overline{Ew}$$

$$\overline{Iwe} = \overline{Ioc} \overline{1} (T6 + T5) \overline{Ew}$$

3.221 Interlace Active

$$sIw = \text{Pot } 1 (T3 - T0) \overline{Ew} \overline{Iw}$$

$$rIw = \overline{Iwf} \overline{T8} (T7 - T0)$$

$$+ (\overline{Wc} + \overline{Iwc} + \overline{Ws} \overline{C19} \overline{C20} \overline{C21} \overline{C22} \overline{C23}) \\ (T3 - T0) \overline{Iw}$$

3.222 Zero Count

$$\begin{aligned} \overline{Iwf} &= \overline{Wc0} \overline{Wc1} \overline{Wc2} \overline{Wc3} \overline{Wc4} \overline{Wc5} \overline{Wc6} \overline{Wc7} \\ &\quad \overline{Wc8} \overline{Wc9} \overline{Wc10} \overline{Wc11} \overline{Wc12} \overline{Wc13} \overline{Wc14} \\ &\quad \overline{Ew} \end{aligned}$$

3.223 Interlace Count Trigger

$$\overline{Iwa} = \overline{Rwx} \overline{Tsm}$$

3.224 Interlace Counter Clock Enables

$$\text{Computer Clock Enable: } \overline{Ew} \overline{Iw}$$

$$\text{Counter Clock Enable: } \overline{Ew}$$

3.225 Extend Operations

$$sIwg = \overline{Iwe} \overline{C12}$$

$$rIwg = \overline{Iwc} + \overline{Wsc} \overline{T8} \overline{W10} \overline{W11} \overline{W12} \overline{W13} \overline{W14} \overline{Iwg}$$

3.226 Channel Command Interrupt Enables

$$\begin{aligned} sIw_j &= Iwe\ C13 && (Eor) \\ rIw_j &= Iwc + I2w\ T8\ Iw_j \\ sIwk &= Iwe\ C14 && (Iwf) \\ rIwk &= Iwc + I1w\ T8\ Iwk \end{aligned}$$

3.227 Channel Command Register

$$\begin{aligned} sIwh &= Iwe\ C15 \\ rIwh &= Iwc \\ sIwi &= Iwe\ C16 \\ rIwi &= Iwc \end{aligned}$$

3.228 Word Counter

$$\begin{aligned} sWc14 &= Iwc + \overline{Ew}\ \overline{Wc14}\ Iwa\ \overline{Ew} \\ rWc14 &= Iwp\ C9 + \overline{Ew}\ \overline{Wc14}\ Iwa\ \overline{Ew} \\ sWc13 &= Iwc + \overline{Ew}\ \overline{Wc13}\ \overline{Wc14}\ \overline{Ew} \\ rWc13 &= Iwp\ C8 + \overline{Ew}\ \overline{Wc13}\ \overline{Wc14}\ \overline{Ew} \\ sWc12 &= Iwc + \overline{Ew}\ \overline{Wc12}\ \overline{Wc13}\ \overline{Wc14}\ \overline{Ew} \\ rWc12 &= Iwp\ C7 + \overline{Ew}\ \overline{Wc12}\ \overline{Wc13}\ \overline{Wc14}\ \overline{Ew} \\ sWc11 &= Iwc + \overline{Ew}\ \overline{Wc11}\ \overline{Wc12}\ \overline{Wc13}\ \overline{Wc14}\ \overline{Ew} \\ rWc11 &= Iwp\ C6 + \overline{Ew}\ \overline{Wc11}\ \overline{Wc12}\ \overline{Wc13}\ \overline{Wc14}\ \overline{Ew} \\ sWc10 &= Iwc + \overline{Ew}\ \overline{Wc10}\ \overline{Wc11}\ \overline{Ew} \\ rWc10 &= Iwp\ C5 + \overline{Ew}\ \overline{Wc10}\ \overline{Wc11}\ \overline{Ew} \\ sWc9 &= Iwc + \overline{Ew}\ \overline{Wc9}\ \overline{Wc10}\ \overline{Wc11}\ \overline{Ew} \\ rWc9 &= Iwp\ C4 + \overline{Ew}\ \overline{Wc9}\ \overline{Wc10}\ \overline{Wc11}\ \overline{Ew} \\ sWc8 &= Iwc + \overline{Ew}\ \overline{Wc8}\ \overline{Wc9}\ \overline{Wc10}\ \overline{Wc11}\ \overline{Ew} \\ rWc8 &= Iwp\ C3 + \overline{Ew}\ \overline{Wc8}\ \overline{Wc9}\ \overline{Wc10}\ \overline{Wc11}\ \overline{Ew} \\ sWc7 &= Iwc + \overline{Ew}\ \overline{Wc7}\ \overline{Wc8}\ \overline{Ew} \\ rWc7 &= Iwp\ C2 + \overline{Ew}\ \overline{Wc7}\ \overline{Wc8}\ \overline{Ew} \\ sWc6 &= Iwc + \overline{Ew}\ \overline{Wc6}\ \overline{Wc7}\ \overline{Wc8}\ \overline{Ew} \\ rWc6 &= Iwp\ C1 + \overline{Ew}\ \overline{Wc6}\ \overline{Wc7}\ \overline{Wc8}\ \overline{Ew} \\ sWc5 &= Iwc + \overline{Ew}\ \overline{Wc5}\ \overline{Wc6}\ \overline{Wc7}\ \overline{Wc8}\ \overline{Ew} \\ rWc5 &= Iwp\ C0 + \overline{Ew}\ \overline{Wc5}\ \overline{Wc6}\ \overline{Wc7}\ \overline{Wc8}\ \overline{Ew} \end{aligned}$$

$$\begin{aligned} sWc4 &= Iwc + \overline{Ew}\ \overline{Wc4}\ \overline{Wc5}\ \overline{Ew} \\ rWc4 &= Iwe\ C23 + \overline{Ew}\ \overline{Wc4}\ \overline{Wc5}\ \overline{Ew} \\ sWc3 &= Iwc + \overline{Ew}\ \overline{Wc3}\ \overline{Wc4}\ \overline{Wc5}\ \overline{Ew} \\ rWc3 &= Iwe\ C22 + \overline{Ew}\ \overline{Wc3}\ \overline{Wc4}\ \overline{Wc5}\ \overline{Ew} \\ sWc2 &= Iwc + \overline{Ew}\ \overline{Wc2}\ \overline{Wc3}\ \overline{Wc4}\ \overline{Wc5}\ \overline{Ew} \\ rWc2 &= Iwe\ C21 + \overline{Ew}\ \overline{Wc2}\ \overline{Wc3}\ \overline{Wc4}\ \overline{Wc5}\ \overline{Ew} \\ sWc1 &= Iwc + \overline{Ew}\ \overline{Wc1}\ \overline{Wc2}\ \overline{Ew} \\ rWc1 &= Iwe\ C20 + \overline{Ew}\ \overline{Wc1}\ \overline{Wc2}\ \overline{Ew} \\ sWc0 &= Iwc + \overline{Ew}\ \overline{Wc0}\ \overline{Wc1}\ \overline{Wc2}\ \overline{Ew} \\ rWc0 &= Iwe\ C19 \end{aligned}$$

3.229 Address Counter

$$\begin{aligned} sWa14 &= Iwc + \overline{Ew}\ \overline{Wa14}\ Iwa\ \overline{Ew} \\ rWa14 &= Iwp\ C23 + \overline{Ew}\ \overline{Wa14}\ Iwa\ \overline{Ew} \\ sWa13 &= Iwc + \overline{Ew}\ \overline{Wa13}\ \overline{Wa14}\ Iwa\ \overline{Ew} \\ rWa13 &= Iwp\ C22 + \overline{Ew}\ \overline{Wa13}\ \overline{Wa14}\ Iwa\ \overline{Ew} \\ sWa12 &= Iwc + \overline{Ew}\ \overline{Wa12}\ \overline{Wa13}\ \overline{Ew} \\ rWa12 &= Iwp\ C21 + \overline{Ew}\ \overline{Wa12}\ \overline{Wa13}\ \overline{Ew} \\ sWa11 &= Iwc + \overline{Ew}\ \overline{Wa11}\ \overline{Wa12}\ \overline{Wa13}\ \overline{Ew} \\ rWa11 &= Iwp\ C20 + \overline{Ew}\ \overline{Wa11}\ \overline{Wa12}\ \overline{Wa13}\ \overline{Ew} \\ sWa10 &= Iwc + \overline{Ew}\ \overline{Wa10}\ \overline{Wa11}\ \overline{Wa12}\ \overline{Wa13}\ \overline{Ew} \\ rWa10 &= Iwp\ C19 + \overline{Ew}\ \overline{Wa10}\ \overline{Wa11}\ \overline{Wa12}\ \overline{Wa13}\ \overline{Ew} \\ sWa9 &= Iwp\ C18 + \overline{Ew}\ \overline{Wa9}\ \overline{Wa10}\ \overline{Ew} \\ rWa9 &= Iwc + \overline{Ew}\ \overline{Wa9}\ \overline{Wa10}\ \overline{Ew} \\ sWa8 &= Iwp\ C17 + \overline{Ew}\ \overline{Wa8}\ \overline{Wa9}\ \overline{Wa10}\ \overline{Ew} \\ rWa8 &= Iwc + \overline{Ew}\ \overline{Wa8}\ \overline{Wa9}\ \overline{Wa10}\ \overline{Ew} \\ sWa7 &= Iwp\ C16 + \overline{Ew}\ \overline{Wa7}\ \overline{Wa8}\ \overline{Wa9}\ \overline{Wa10}\ \overline{Ew} \\ rWa7 &= Iwc + \overline{Ew}\ \overline{Wa7}\ \overline{Wa8}\ \overline{Wa9}\ \overline{Wa10}\ \overline{Ew} \\ sWa6 &= Iwp\ C15 + \overline{Ew}\ \overline{Wa6}\ \overline{Wa7}\ \overline{Ew} \\ rWa6 &= Iwc + \overline{Ew}\ \overline{Wa6}\ \overline{Wa7}\ \overline{Ew} \\ sWa5 &= Iwp\ C14 + \overline{Ew}\ \overline{Wa5}\ \overline{Wa6}\ \overline{Wa7}\ \overline{Ew} \\ rWa5 &= Iwc + \overline{Ew}\ \overline{Wa5}\ \overline{Wa6}\ \overline{Wa7}\ \overline{Ew} \end{aligned}$$

$$sWa4 = Iwp C13 + \overline{Ew} \overline{Wa4} \overline{Wa5} \overline{Wa6} \overline{Wa7} \overline{Ew}$$

$$rWa4 = Iwc + \overline{Ew} \overline{Wa4} \overline{Wa5} \overline{Wa6} \overline{Wa7} \overline{Ew}$$

$$sWa3 = Iwp C12 + \overline{Ew} \overline{Wa3} \overline{Wa4} \overline{Ew}$$

$$rWa3 = Iwc + \overline{Ew} \overline{Wa3} \overline{Wa4} \overline{Ew}$$

$$sWa2 = Iwp C11 + \overline{Ew} \overline{Wa2} \overline{Wa3} \overline{Wa4} \overline{Ew}$$

$$rWa2 = Iwc + \overline{Ew} \overline{Wa2} \overline{Wa3} \overline{Wa4} \overline{Ew}$$

$$sWa1 = Iwp C10 + \overline{Ew} \overline{Wa1} \overline{Wa2} \overline{Wa3} \overline{Wa4} \overline{Ew}$$

$$rWa1 = Iwc + \overline{Ew} \overline{Wa1} \overline{Wa2} \overline{Wa3} \overline{Wa4} \overline{Ew}$$

$$sWa0 = Iwe C18 + \overline{Ew} \overline{Wa0} \overline{Wa1} \overline{Ew}$$

$$rWa0 = Iwc$$

$$Zw7 = \textcircled{Zw7}$$

$$Zw8 = \textcircled{Zw8}$$

$$Zw9 = \textcircled{Zw9}$$

$$Zw10 = \textcircled{Zw10}$$

$$Zw11 = \textcircled{Zw11}$$

$$Zw12 = \textcircled{Zw12}$$

$$Zw13 = \textcircled{Zw13}$$

$$Zw14 = \textcircled{Zw14}$$

$$Zw15 = \textcircled{Zw15}$$

$$Zw16 = \textcircled{Zw16}$$

$$Zw17 = \textcircled{Zw17}$$

$$Zw18 = \textcircled{Zw18}$$

$$Zw19 = \textcircled{Zw19}$$

$$Zw20 = \textcircled{Zw20}$$

$$Zw21 = \textcircled{Zw21}$$

$$Zw22 = \textcircled{Zw22}$$

$$Zw23 = \textcircled{Zw23}$$

$$Whs = \textcircled{Whs}$$

$$Wes = \textcircled{Wes}$$

$$Mtgw = \textcircled{Mtgw}$$

$$Npw = \textcircled{Npw}$$

$$Ecw = \textcircled{Ecw}$$

$$Wx12 = \textcircled{Wx12} \quad (93200 \text{ and } 93221 \text{ only})$$

$$Wx24 = \textcircled{Wx24} \quad (93200 \text{ and } 93221 \text{ only})$$

3.230 PIN Address Counter

$$sWpa = Ioc \overline{C17} \overline{C13} \overline{C19} \overline{C20} \overline{C21} \overline{C22} \overline{C23} T0$$

$$rWpa = Pin T0 + St$$

3.231 Skip Gate:

$$\begin{aligned} Skr &= \overline{C1} \overline{C17} \overline{C9} \overline{C10} \overline{C11} \overline{C19} \overline{C20} \overline{C21} \overline{C22} \\ &\quad \overline{C23} C15 Wsc* \\ &+ \overline{C1} \overline{C17} \overline{C9} \overline{C10} \overline{C11} \overline{C19} \overline{C20} \overline{C21} \overline{C22} \\ &\quad \overline{C23} C12 \overline{W10} \overline{W11} \overline{W12} \overline{W13} \overline{W14}* \\ &+ \overline{C1} \overline{C17} \overline{C9} \overline{C10} \overline{C11} \overline{C19} \overline{C20} \overline{C21} \overline{C22} \\ &\quad \overline{C23} C13 Iwf* \\ &+ \overline{C1} \overline{C17} \overline{C9} \overline{C10} \overline{C11} \overline{C19} \overline{C20} \overline{C21} \overline{C22} \\ &\quad \overline{C23} C14 \overline{We}* \\ &+ C10 \overline{C11} C14 \overline{W10} \overline{W11} \overline{W12} \overline{W13} \overline{W14} C1* \\ &+ C10 \overline{C11} C20 We \overline{C1}* \\ &+ \dots \end{aligned}$$

* $\overline{C1}$ becomes C1 for TMCC-C

3.232 Input/Output Signals Received

$$Zwp = \textcircled{Zwp}$$

$$Zw1 = \textcircled{Zw1}$$

$$Zw2 = \textcircled{Zw2}$$

$$Zw3 = \textcircled{Zw3}$$

$$Zw4 = \textcircled{Zw4}$$

$$Zw5 = \textcircled{Zw5}$$

$$Zw6 = \textcircled{Zw6}$$

3.233 Input/Output Signals Generated

$$\textcircled{Rwp} = Rwp$$

$$\textcircled{Rw1} = Rw1$$

$$\textcircled{Rw2} = Rw2$$

$$\textcircled{Rw3} = Rw3$$

$$\textcircled{Rw4} = Rw4$$

$$\textcircled{Rw5} = Rw5$$

$$\textcircled{Rw6} = Rw6$$

$$\textcircled{Rw7} = Rw7$$

$$\textcircled{Rw8} = Rw8$$

$$\textcircled{Rw9} = Rw9$$

$$\textcircled{Rw10} = Rw10$$

$$\textcircled{Rw11} = Rw11$$

- $\textcircled{\text{Rw12}}$ = Rw12
- $\textcircled{\text{Rw13}}$ = Rw13
- $\textcircled{\text{Rw14}}$ = Rw14
- $\textcircled{\text{Rw15}}$ = Rw15
- $\textcircled{\text{Rw16}}$ = Rw16
- $\textcircled{\text{Rw17}}$ = Rw17
- $\textcircled{\text{Rw18}}$ = Rw18
- $\textcircled{\text{Rw19}}$ = Rw19
- $\textcircled{\text{Rw20}}$ = Rw20
- $\textcircled{\text{Rw21}}$ = Rw21
- $\textcircled{\text{Rw22}}$ = Rw22
- $\textcircled{\text{Rw23}}$ = Rw23
- $\textcircled{\text{Rw24}}$ = Rw24
- $\textcircled{\text{W9}}$ = W9
- $\textcircled{\text{W10}}$ = W10
- $\textcircled{\text{W11}}$ = W11
- $\textcircled{\text{W12}}$ = W12
- $\textcircled{\text{W13}}$ = W13
- $\textcircled{\text{W0}}$ = W0
- $\textcircled{\text{W5}}$ = W5
- $\textcircled{\text{W6}}$ = W6
- $\textcircled{\text{Iw}}$ = Iw

3.234 LOGIC EQUATIONS FOR Y BUFFER

3.235 Unit Address Register

- $sY14 = Y_s C23$
- $rY14 = Y_c$
- $sY13 = Y_s C22$
- $rY13 = Y_c$
- $sY12 = Y_s C21$
- $rY12 = Y_c$
- $sY11 = Y_s C20$
- $rY11 = Y_c$
- $sY10 = Y_s C19 + (\text{Ioc } C12 C17 \overline{C19} \overline{C20} \overline{C21} \overline{C22} \overline{C23}) \overline{Y9} \overline{Y10}$
- $rY10 = Y_c$

3.236 Input/Output

- $sY9 = Y_s C18$
- $rY9 = Y_c$

3.237 Clear and Set Signals

- $Y_c = \text{Buc } C17 (T6 + T5) + Y_h \overline{Y_f} (T3 - T0) + S_t$
- $Y_s = \text{Buc } C17 (T3 - T0)$

3.238 Clock Counter

- $sY6 = \overline{Y5} E_{cy} T8 \overline{Y10} \overline{Y11} \overline{Y12} \overline{Y13} \overline{Y14}$
- $rY6 = Y5 T0 + Y_c$
- $sY5 = \overline{Y5} Y6 \overline{E_{cy}} T0 + Y_s C13 C18 T0 \overline{Y10} \overline{Y11} \overline{Y12} \overline{Y13} \overline{Y14}$
- $rY5 = Y4 T0 + Y_c$
- $sY4 = Y5 Y_f T8 \overline{Y_g} + Y_s T0 + Y_h Y_f T8 + I_{yg} \overline{Y9} \overline{I_{yi}} Y5 I_{yf} T8 + I_{yg} Y_g Y_f Y0 \overline{Y_{ev}} I_y T8 \overline{Y7} \overline{Y9} \overline{Y10} \overline{Y11} \overline{Y_h}$
- $rY4 = Y4 T0 + Y4 T8$

3.239 Character Counter

- $sY8 = Y_s C16 + Y7 \overline{Y8} Y4 T0 + Y_{xx} Y_{n2} (\overline{T7} - \overline{T0}) \overline{Y4} + \overline{Y7} \overline{Y9} Y10 Y11 \overline{Y_h}$
- $rY8 = Y_c + Y8 Y4 T0$
- $sY7 = Y_s C15 + Y_{xx} Y_{n1} (\overline{T7} - \overline{T0}) Y4$
- $rY7 = Y_c + Y7 \overline{Y8} Y4 T0$

3.240 Character Counter Even

- $Y_{ev} = Y8 Y_{n2} Y7 Y_{n1} + \overline{Y8} \overline{Y_{n2}} Y7 Y_{n1} + \overline{Y8} \overline{Y_{n2}} \overline{Y7} \overline{Y_{n1}} + Y8 Y_{n2} \overline{Y7} \overline{Y_{n1}}$

3.241 Halt Interlock

$$\begin{aligned}
 sY0 &= \overline{Y9} Y6 \overline{Y8} Ecy \\
 &+ Ys C18 Y9 \\
 rY0 &= Yc \\
 &+ (Ioc C12 C17 \overline{C19} \overline{C20} \overline{C21} \overline{C22} \overline{C23}) \\
 &Y9 T0 \\
 &+ Y9 Iy Iyf (\overline{Iyg} + \overline{Iyh} + \overline{Iyi}) \overline{T7 - T0}
 \end{aligned}$$

3.242 Computer Interlock

$$\begin{aligned}
 sYf &= Yc \overline{Yh} \\
 &+ Iya \overline{Mit} T0 \\
 &+ Rx T0 \overline{Pwy} * (\overline{Pwy} = \overline{05} \text{ for } 92210) \\
 rYf &= \overline{Y7} \overline{Y8} Y4 (T6 + T5) \\
 &+ Ys C18 \\
 &+ \overline{Y9} Y10 Y11 Y0 Mtgy \overline{Y7} (T6 + T5) \overline{Yh}
 \end{aligned}$$

*Rx is always false for TMCC-D

3.243 End-of-Record Detector

$$\begin{aligned}
 sYg &= Mtgy T0 Iyg Y11 (Y0 \overline{Y9} + \overline{Y0} Y5 \overline{Y6} Y9) \\
 &+ Yhs \overline{T7 - T0} \overline{Y10} \overline{Y11} \overline{Y12} \overline{Y13} \overline{Y14} \\
 &+ \overline{Y9} \overline{Y10} \overline{Y11} Y12 Y13 (\overline{Ry1} \overline{Ry2} \overline{Ry3} \overline{Ry4} \\
 &\overline{Ry5} \overline{Ry6} \overline{Ryp}) Y5 \overline{T7 - T0} \overline{Y10} \overline{Y11} \overline{Y12} \\
 &\overline{Y13} \overline{Y14} \\
 rYg &= Yc \\
 &+ Yg \overline{Iyg} Y11 T0
 \end{aligned}$$

3.244 Halt Detector

$$\begin{aligned}
 sYh &= Yg \overline{Iyg} T8 \\
 &+ Yhs Y11 T8 \\
 &+ Yg \overline{Y9} \overline{Y11} \overline{Iyh} (Iyf + Yev Yf) T8 Iyg \\
 &+ Y9 \overline{Y11} \overline{Y0} Y5 \overline{Y6} (\overline{Iyg} + \overline{Iyi} \overline{Iyh}) T8 \\
 &+ \overline{Y9} \overline{Y11} Iyg \overline{Iyh} Iyi Iyf T8 Y0 \\
 &+ Y9 \overline{Iyh} Yhs T8 \\
 rYh &= Yc (T6 + T5) \\
 &+ Yh \overline{Yf} T8
 \end{aligned}$$

3.245 Signal Complete

$$\begin{aligned}
 sYsc &= Yg Iyg Iyh (Yev Yf + Iyf) T8 \\
 &+ Yh \overline{Yf} T8 + St \\
 &+ Ys \overline{C19} \overline{C20} \overline{C21} \overline{C22} \overline{C23} \\
 rYsc &= Ysc T0 \overline{Iyg}
 \end{aligned}$$

3.246 Interrupt Signals

$$\begin{aligned}
 I1y &= \overline{Yf} Y0 \overline{Yh} (En + \textcircled{En}) \overline{Iy} \overline{Ey} \overline{Iyg} \\
 &+ Iyg \overline{Iy} Iyf Iyk \\
 I2y &= (En + \textcircled{En}) \overline{Iyg} Yh \overline{Yf} \\
 &+ Ysc Iyj Iyg
 \end{aligned}$$

3.247 YIM + MIY Interlock =

$$\overline{Yf} (Y0 + \overline{Y9})$$

3.248 Load Buffer from C

$$Yxx = Rx \overline{Pwy} + Ryx Tsr Iy * (\overline{Pwy} = \overline{05} \text{ and } Iy \text{ deleted for } 92210)$$

*Rx is always false for TMCC-D

3.249 Time Share Request

$$Trqy = \overline{Yf} Y0 Yh Iy \overline{Iyf}$$

3.250 Time Share Select

$$\begin{aligned}
 sRyx &= \overline{Tsm} Trqy \overline{Trqx} (\overline{T7 - T0}) \\
 rRyx &= \overline{Tsm} T0
 \end{aligned}$$

3.251 Time Share Priority

$$\begin{aligned}
 Trqx &= \text{Priority Signal from TMCC-C and TMCC-D} \\
 Trqx &= Trq(c) + Trq(d) \\
 Trq &= Trqx + Trqw + Trqy
 \end{aligned}$$

3.252 Y Register

$$\begin{aligned}
 Yw1 &= Y4 Y7 (\overline{T7 - T0}) \\
 &+ \overline{Y4} Yn1 (\overline{T7 - T0}) \\
 &+ \overline{Y4} Yn1 \overline{Yxx}
 \end{aligned}$$

$$+ Y4 Yb1 (T7 - T0) *$$

$$+ \overline{Y4} C21r (T7 - T0) Yxx$$

$$Yw2 = Y4 Y8 (\overline{T7 - T0})$$

$$+ \overline{Y4} Yn2 (\overline{T7 - T0})$$

$$+ \overline{Y4} Yn2 \overline{Yxx}$$

$$+ Y4 Yb2 (T7 - T0) *$$

$$+ \overline{Y4} C23r (T7 - T0) Yxx$$

$$Yw3 = \overline{Y4} Yn3 \overline{Yxx}$$

$$+ Y4 Yb3 *$$

$$+ \overline{Y4} C23r Yxx$$

sYr1,2,3 = Yw1, 2, 3 delayed by 9 pulse times

sYn1,2,3 = Yr1, 2, 3 respectively

$$* \left. \begin{aligned} Yb1 &= \overline{Yx12} \overline{Yx24} Ry4 + Yx12 Ry10 \\ &\quad + Yx24 Ry22 \\ Yb2 &= \overline{Yx12} \overline{Yx24} Ry5 + Yx12 Ry11 \\ &\quad + Yx24 Ry23 \\ Yb3 &= \overline{Yx12} \overline{Yx24} Ry6 + Yx12 Ry12 \\ &\quad + Yx24 Ry24 \end{aligned} \right\} \text{ for 93221}$$

Yb1,2,3 equals Ry4,5,6 respectively for 92210

Yb1,2,3 equals Ry10, 11, 12 respectively for 92211

Yb1,2,3 equals Ry22, 23, 24 respectively for 92212

3.253 Character Buffer Extended to 12 Bits

$$sRy7 = Y4 \overline{Ry7} Ry4 + \overline{Y9} Y6 \overline{Y5} Zy7$$

$$rRy7 = Y4 Ry7 \overline{Ry4} + \overline{Y9} \overline{Y6} \overline{Y5} \overline{Y4}$$

$$sRy8 = Y4 \overline{Ry8} Ry5 + \overline{Y9} Y6 \overline{Y5} Zy8$$

$$rRy8 = Y4 Ry8 \overline{Ry5} + \overline{Y9} \overline{Y6} \overline{Y5} \overline{Y4}$$

$$sRy9 = Y4 \overline{Ry9} Ry6 + \overline{Y9} Y6 \overline{Y5} Zy9$$

$$rRy9 = Y4 Ry9 \overline{Ry6} + \overline{Y9} \overline{Y6} \overline{Y5} \overline{Y4}$$

$$sRy10 = Y4 \overline{Ry10} Ry7 + \overline{Y9} Y6 \overline{Y5} Zy10$$

$$rRy10 = Y4 Ry10 \overline{Ry7} + \overline{Y9} \overline{Y6} \overline{Y5} \overline{Y4}$$

$$sRy11 = Y4 \overline{Ry11} Ry8 + \overline{Y9} Y6 \overline{Y5} Zy11$$

$$rRy11 = Y4 Ry11 \overline{Ry8} + \overline{Y9} \overline{Y6} \overline{Y5} \overline{Y4}$$

$$sRy12 = Y4 \overline{Ry12} Ry9 + \overline{Y9} Y6 \overline{Y5} Zy12$$

$$rRy12 = Y4 Ry12 \overline{Ry9} + \overline{Y9} \overline{Y6} \overline{Y5} \overline{Y4}$$

3.254 Character Buffer *

$$sRy1 = Y4 \overline{Yxx} (T7 - T0) Yn1 + \overline{Y9} Y6 \overline{Y5} Zy1 + Y4 Yxx C21r$$

$$rRy1 = Y4 \overline{Yxx} \overline{Yn1} + \overline{Y9} \overline{Y6} \overline{Y5} \overline{Y4} + Y4 Yxx \overline{C21r}$$

$$sRy2 = Y4 \overline{Yxx} (T7 - T0) Yn2 + \overline{Y9} Y6 \overline{Y5} Zy2 + Y4 Yxx C22r$$

$$rRy2 = Y4 \overline{Yxx} \overline{Yn2} + \overline{Y9} \overline{Y6} \overline{Y5} \overline{Y4} + Y4 Yxx \overline{C22r}$$

$$sRy3 = Y4 \overline{Yxx} (T7 - T0) Yn3 + \overline{Y9} Y6 \overline{Y5} Zy3 + Y4 Yxx C23r$$

$$rRy3 = Y4 \overline{Yxx} \overline{Yn3} + \overline{Y9} \overline{Y6} \overline{Y5} \overline{Y4} + Y4 Yxx \overline{C23r}$$

$$sRy4 = Y4 \overline{Ry4} Ry1 + \overline{Y9} Y6 \overline{Y5} Zy4$$

$$rRy4 = Y4 Ry4 \overline{Ry1} + \overline{Y9} \overline{Y6} \overline{Y5} \overline{Y4}$$

$$sRy5 = Y4 \overline{Ry5} Ry2 + \overline{Y9} Y6 \overline{Y5} Zy5$$

$$rRy5 = Y4 Ry5 \overline{Ry2} + \overline{Y9} \overline{Y6} \overline{Y5} \overline{Y4}$$

$$sRy6 = Y4 \overline{Ry6} Ry3 + \overline{Y9} Y6 \overline{Y5} Zy6$$

$$rRy6 = Y4 Ry6 \overline{Ry3} + \overline{Y9} \overline{Y6} \overline{Y5} \overline{Y4}$$

* For 12-bit extension, add Ry7 through Ry12
For 24-bit extension, add Ry7 through Ry24

Character Size

12 bit character = Yx12 (93221 only)

24 bit character = Yx24 (93221 only)

3.255 Character Buffer Extended to 24 bits

$$sRy13 = Y4 \overline{Ry13} Ry10 + \overline{Y9} Y6 \overline{Y5} Zy13$$

$$rRy13 = Y4 Ry13 \overline{Ry10} + \overline{Y9} \overline{Y6} \overline{Y5} \overline{Y4}$$

$$sRy14 = Y4 \overline{Ry14} Ry11 + \overline{Y9} Y6 \overline{Y5} Zy14$$

$$rRy14 = Y4 Ry14 \overline{Ry11} + \overline{Y9} \overline{Y6} \overline{Y5} \overline{Y4}$$

$$\begin{aligned}
 sRy15 &= Y4 \overline{Ry15} Ry12 + \overline{Y9} Y6 \overline{Y5} Zy15 \\
 rRy15 &= Y4 Ry15 \overline{Ry12} + \overline{Y9} \overline{Y6} \overline{Y5} \overline{Y4} \\
 sRy16 &= Y4 \overline{Ry16} Ry13 + \overline{Y9} Y6 \overline{Y5} Zy16 \\
 rRy16 &= Y4 Ry16 \overline{Ry13} + \overline{Y9} \overline{Y6} \overline{Y5} \overline{Y4} \\
 sRy17 &= Y4 \overline{Ry17} Ry14 + \overline{Y9} Y6 \overline{Y5} Zy17 \\
 rRy17 &= Y4 Ry17 \overline{Ry14} + \overline{Y9} \overline{Y6} \overline{Y5} \overline{Y4} \\
 sRy18 &= Y4 \overline{Ry18} Ry15 + \overline{Y9} Y6 \overline{Y5} Zy18 \\
 rRy18 &= Y4 Ry18 \overline{Ry15} + \overline{Y9} \overline{Y6} \overline{Y5} \overline{Y4} \\
 sRy19 &= Y4 \overline{Ry19} Ry16 + \overline{Y9} Y6 \overline{Y5} Zy19 \\
 rRy19 &= Y4 Ry19 \overline{Ry16} + \overline{Y9} \overline{Y6} \overline{Y5} \overline{Y4} \\
 sRy20 &= Y4 \overline{Ry20} Ry17 + \overline{Y9} Y6 \overline{Y5} Zy20 \\
 rRy20 &= Y4 Ry20 \overline{Ry17} + \overline{Y9} \overline{Y6} \overline{Y5} \overline{Y4} \\
 sRy21 &= Y4 \overline{Ry21} Ry18 + \overline{Y9} Y6 \overline{Y5} Zy21 \\
 rRy21 &= Y4 Ry21 \overline{Ry18} + \overline{Y9} \overline{Y6} \overline{Y5} \overline{Y4} \\
 sRy22 &= Y4 \overline{Ry22} Ry19 + \overline{Y9} Y6 \overline{Y5} Zy22 \\
 rRy22 &= Y4 Ry22 \overline{Ry19} + \overline{Y9} \overline{Y6} \overline{Y5} \overline{Y4} \\
 sRy23 &= Y4 \overline{Ry23} Ry20 + \overline{Y9} Y6 \overline{Y5} Zy23 \\
 rRy23 &= Y4 Ry23 \overline{Ry20} + \overline{Y9} \overline{Y6} \overline{Y5} \overline{Y4} \\
 sRy24 &= Y4 \overline{Ry24} Ry21 + \overline{Y9} Y6 \overline{Y5} Zy24 \\
 rRy24 &= Y4 Ry24 \overline{Ry21} + \overline{Y9} \overline{Y6} \overline{Y5} \overline{Y4}
 \end{aligned}$$

3.256 Parity Flip-Flop

$$\begin{aligned}
 sRyp &= \overline{Y9} Y4 \overline{Ryp} (Yb1 \oplus Yb2 \oplus Yb3) (T7 - T0) Qy1 * \\
 &+ Y9 Y4 \overline{Ryp} \overline{Yxx} (Yn1 \oplus Yn2 \oplus Yn3) Qy2 \\
 &\quad (T7 - T0) * \\
 &+ Y9 Y4 \overline{Ryp} Yxx Cpr Qy2 (T7 - T0) * \\
 &+ \overline{Y9} Y6 \overline{Y5} Zyp \\
 &+ Yf Y5 T8 \overline{Ryp} \\
 rRyp &= \overline{Y9} Y4 Ryp (Yb1 \oplus Yb2 \oplus Yb3) (T7 - T0) \\
 &\quad Qy1 * \\
 &+ Y9 Y4 Ryp \overline{Yxx} (Yn1 \oplus Yn2 \oplus Yn3) \\
 &\quad Qy2 (T7 - T0) * \\
 &+ Y9 Y4 Ryp Yxx Cpr Qy2 (T7 - T0) * \\
 &+ \overline{Y9} \overline{Y6} \overline{Y5} \overline{Y4}
 \end{aligned}$$

$$\begin{aligned}
 &+ Yf Y5 T8 \overline{Y9} Ryp \\
 &+ Yc \\
 * Qy1 &= \overline{Yx12} \overline{Qr4} + \overline{Yx12} \overline{Yx24} \overline{Qr1} \overline{Qr4} \text{ for 93221} \\
 Qy2 &= Yx12 Qr4 + \overline{Yx12} \overline{Yx24} \overline{Qr2} \overline{Qr4} \text{ for 93221} \\
 Qy1 &= Qr1 Qr4 \text{ for 92210} \\
 Qy1 &= Qr4 \text{ for 92211} \\
 Qr1 &\text{ is deleted for 92212} \\
 Qy2 &= Qr2 \overline{Qr4} \text{ for 92210} \\
 Qy2 &= \overline{Qr4} \text{ for 92211} \\
 Qy2 &\text{ is deleted for 92212}
 \end{aligned}$$

3.257 Error Detector

$$\begin{aligned}
 sYe &= \overline{Y9} \overline{Y6} \overline{Y5} \overline{Y4} Ryp \overline{Yg} Npy (\overline{Iyg} + Iyi + \overline{Iyf}) \\
 &+ Y0 \overline{Y6} Y5 Ecy T8 \\
 &+ Yes \\
 rYe &= Yc \overline{Yh}
 \end{aligned}$$

3.258 Interlace Prepare

$$\begin{aligned}
 sEy &= Iyc \overline{Ey} (T3 - T0) \\
 rEy &= Yc T0 \\
 &+ Pot 1 (T3 - T0) Ey
 \end{aligned}$$

3.259 Interlace Clear

$$Iyc = Eom C9 \overline{C10} \overline{C1} C17 (T3 - T0) *$$

* $\overline{C1}$ becomes C1 for TMCC-D

3.260 Interlace Load

$$\begin{aligned}
 Iyp &= Pot 1 (T6 + T5) Ey \\
 Iye &= Ioc 1 Ey (T6 + T5)
 \end{aligned}$$

3.261 Interlace Active

$$\begin{aligned}
 sIy &= Pot 1 (T3 - T0) Ey \overline{Iy} \\
 rIy &= Iyf T8 \\
 &+ (Yc + Iyc + Ys \overline{C19} \overline{C20} \overline{C21} \overline{C22} \overline{C23}) \\
 &\quad (T3 - T0) Iy
 \end{aligned}$$

3.262 Zero Count

$$I_{yf} = Y_{c0} Y_{c1} Y_{c2} Y_{c3} Y_{c4} Y_{c5} Y_{c6} Y_{c7} Y_{c8} Y_{c9} Y_{c10} Y_{c11} Y_{c12} Y_{c13} Y_{c14} \bar{E}_y$$

3.263 Interlace Count Trigger

$$I_{ya} = R_{yx} T_{sm}$$

3.264 Interlace Counter Clock Enables

Computer Clock Enable: $E_y \bar{I}_y$

Counter Clock Enable: \bar{E}_y

3.265 Extend Operations

$$sI_{yg} = I_{ye} C_{12}$$

$$rI_{yg} = I_{yc} + Y_{sc} T_8 \bar{Y}_{10} \bar{Y}_{11} \bar{Y}_{12} \bar{Y}_{13} \bar{Y}_{14} I_{yg}$$

3.266 Channel Command Interrupt Enables

$$sI_{yj} = I_{ye} C_{13}$$

$$rI_{yj} = I_{yc} + I_{2y} T_8 I_{yj}$$

$$sI_{yk} = I_{ye} C_{14}$$

$$rI_{yk} = I_{yc} + I_{1y} T_8 I_{yk}$$

3.267 Channel Command Register

$$sI_{yh} = I_{ye} C_{15}$$

$$rI_{yh} = I_{yc}$$

$$sI_{yi} = I_{ye} C_{16}$$

$$rI_{yi} = I_{yc}$$

3.268 Word Counter

$$sY_{c14} = I_{yc} + \bar{E}_y \bar{Y}_{c14} I_{ya} \bar{E}_y$$

$$rY_{c14} = I_{yp} C_9 + \bar{E}_y Y_{c14} I_{ya} \bar{E}_y$$

$$sY_{c13} = I_{yc} + \bar{E}_y \bar{Y}_{c13} Y_{c14} \bar{E}_y$$

$$rY_{c13} = I_{yp} C_8 + \bar{E}_y Y_{c13} Y_{c14} \bar{E}_y$$

$$sY_{c12} = I_{yc} + \bar{E}_y \bar{Y}_{c12} Y_{c13} Y_{c14} \bar{E}_y$$

$$rY_{c12} = I_{yp} C_7 + \bar{E}_y Y_{c12} Y_{c13} Y_{c14} \bar{E}_y$$

$$sY_{c11} = I_{yc} + \bar{E}_y \bar{Y}_{c11} Y_{c12} Y_{c13} Y_{c14} \bar{E}_y$$

$$rY_{c11} = I_{yp} C_6 + \bar{E}_y Y_{c11} Y_{c12} Y_{c13} Y_{c14} \bar{E}_y$$

$$sY_{c10} = I_{yc} + \bar{E}_y \bar{Y}_{c10} Y_{c11} \bar{E}_y$$

$$rY_{c10} = I_{yp} C_5 + \bar{E}_y Y_{c10} Y_{c11} \bar{E}_y$$

$$sY_{c9} = I_{yc} + \bar{E}_y \bar{Y}_{c9} Y_{c10} Y_{c11} \bar{E}_y$$

$$rY_{c9} = I_{yp} C_4 + \bar{E}_y Y_{c9} Y_{c10} Y_{c11} \bar{E}_y$$

$$sY_{c8} = I_{yc} + \bar{E}_y \bar{Y}_{c8} Y_{c9} Y_{c10} Y_{c11} \bar{E}_y$$

$$rY_{c8} = I_{yp} C_3 + \bar{E}_y Y_{c8} Y_{c9} Y_{c10} Y_{c11} \bar{E}_y$$

$$sY_{c7} = I_{yc} + \bar{E}_y \bar{Y}_{c7} Y_{c8} \bar{E}_y$$

$$rY_{c7} = I_{yp} C_2 + \bar{E}_y Y_{c7} Y_{c8} \bar{E}_y$$

$$sY_{c6} = I_{yc} + \bar{E}_y \bar{Y}_{c6} Y_{c7} Y_{c8} \bar{E}_y$$

$$rY_{c6} = I_{yp} C_1 + \bar{E}_y Y_{c6} Y_{c7} Y_{c8} \bar{E}_y$$

$$sY_{c5} = I_{yc} + \bar{E}_y \bar{Y}_{c5} Y_{c6} Y_{c7} Y_{c8} \bar{E}_y$$

$$rY_{c5} = I_{yp} C_0 + \bar{E}_y Y_{c5} Y_{c6} Y_{c7} Y_{c8} \bar{E}_y$$

$$sY_{c4} = I_{yc} + \bar{E}_y \bar{Y}_{c4} Y_{c5} \bar{E}_y$$

$$rY_{c4} = I_{ye} C_{23} + \bar{E}_y Y_{c4} Y_{c5} \bar{E}_y$$

$$sY_{c3} = I_{yc} + \bar{E}_y \bar{Y}_{c3} Y_{c4} Y_{c5} \bar{E}_y$$

$$rY_{c3} = I_{ye} C_{22} + \bar{E}_y Y_{c3} Y_{c4} Y_{c5} \bar{E}_y$$

$$sY_{c2} = I_{yc} + \bar{E}_y \bar{Y}_{c2} Y_{c3} Y_{c4} Y_{c5} \bar{E}_y$$

$$rY_{c2} = I_{ye} C_{21} + \bar{E}_y Y_{c2} Y_{c3} Y_{c4} Y_{c5} \bar{E}_y$$

$$sY_{c1} = I_{yc} + \bar{E}_y \bar{Y}_{c1} Y_{c2} \bar{E}_y$$

$$rY_{c1} = I_{ye} C_{20} + \bar{E}_y Y_{c1} Y_{c2} \bar{E}_y$$

$$sY_{c0} = I_{yc} + \bar{E}_y \bar{Y}_{c0} Y_{c1} Y_{c2} \bar{E}_y$$

$$rY_{c0} = I_{ye} C_{19}$$

3.269 Address Counter

$$sY_{a14} = I_{yc} + \bar{E}_y \bar{Y}_{a14} I_{ya} \bar{E}_y$$

$$rY_{a14} = I_{yp} C_{23} + \bar{E}_y Y_{a14} I_{ya} \bar{E}_y$$

$$sY_{a13} = I_{yc} + \bar{E}_y \bar{Y}_{a13} Y_{a14} I_{ya} \bar{E}_y$$

$$rY_{a13} = I_{yp} C_{22} + \bar{E}_y Y_{a13} Y_{a14} I_{ya} \bar{E}_y$$

$$sY_{a12} = I_{yc} + \bar{E}_y \bar{Y}_{a12} Y_{a13} \bar{E}_y$$

$$rY_{a12} = I_{yp} C_{21} + \bar{E}_y Y_{a12} Y_{a13} \bar{E}_y$$

$$sY_{a11} = I_{yc} + \bar{E}_y \bar{Y}_{a11} Y_{a12} Y_{a13} \bar{E}_y$$

$$rY_{a11} = I_{yp} C_{20} + \bar{E}_y Y_{a11} Y_{a12} Y_{a13} \bar{E}_y$$

$$sY_{a10} = I_{yc} + \bar{E}_y \bar{Y}_{a10} Y_{a11} Y_{a12} Y_{a13} \bar{E}_y$$

$$rY_{a10} = I_{yp} C_{19} + \bar{E}_y Y_{a10} Y_{a11} Y_{a12} Y_{a13} \bar{E}_y$$

$$sYa9 = Iyp C18 + \overline{E_y} \overline{Ya9} \underline{Ya10 \overline{E_y}} + C10 \overline{C11} C19 \overline{Y_e} \overline{C1}^*$$

$$rYa9 = Iyc + \overline{E_y} \overline{Ya9} \underline{Ya10 \overline{E_y}} + \dots$$

$$sYa8 = Iyp C17 + \overline{E_y} \overline{Ya8} \underline{Ya9 \overline{Ya10 \overline{E_y}}}$$

$$rYa8 = Iyc + \overline{E_y} \overline{Ya8} \underline{Ya9 \overline{Ya10 \overline{E_y}}}$$

* $\overline{C1}$ becomes C1 for TMCC-D

$$sYa7 = Iyp C16 + \overline{E_y} \overline{Ya7} \underline{Ya8 \overline{Ya9 \overline{Ya10 \overline{E_y}}}}$$

$$rYa7 = Iyc + \overline{E_y} \overline{Ya7} \underline{Ya8 \overline{Ya9 \overline{Ya10 \overline{E_y}}}}$$

3.272 Input/Output Signals Received

$$sYa6 = Iyp C15 + \overline{E_y} \overline{Ya6} \underline{Ya7 \overline{E_y}}$$

$$rYa6 = Iyc + \overline{E_y} \overline{Ya6} \underline{Ya7 \overline{E_y}}$$

$$Zyp = \overline{\overline{Zyp}}$$

$$Zy1 = \overline{\overline{Zy1}}$$

$$Zy2 = \overline{\overline{Zy2}}$$

$$sYa5 = Iyp C14 + \overline{E_y} \overline{Ya5} \underline{Ya6 \overline{Ya7 \overline{E_y}}}$$

$$rYa5 = Iyc + \overline{E_y} \overline{Ya5} \underline{Ya6 \overline{Ya7 \overline{E_y}}}$$

$$Zy3 = \overline{\overline{Zy3}}$$

$$Zy4 = \overline{\overline{Zy4}}$$

$$sYa4 = Iyp C13 + \overline{E_y} \overline{Ya4} \underline{Ya5 \overline{Ya6 \overline{Ya7 \overline{E_y}}}}$$

$$rYa4 = Iyc + \overline{E_y} \overline{Ya4} \underline{Ya5 \overline{Ya6 \overline{Ya7 \overline{E_y}}}}$$

$$Zy5 = \overline{\overline{Zy5}}$$

$$Zy6 = \overline{\overline{Zy6}}$$

$$sYa3 = Iyp C12 + \overline{E_y} \overline{Ya3} \underline{Ya4 \overline{E_y}}$$

$$rYa3 = Iyc + \overline{E_y} \overline{Ya3} \underline{Ya4 \overline{E_y}}$$

$$Zy7 = \overline{\overline{Zy7}}$$

$$Zy8 = \overline{\overline{Zy8}}$$

$$sYa2 = Iyp C11 + \overline{E_y} \overline{Ya2} \underline{Ya3 \overline{Ya4 \overline{E_y}}}$$

$$rYa2 = Iyc + \overline{E_y} \overline{Ya2} \underline{Ya3 \overline{Ya4 \overline{E_y}}}$$

$$Zy9 = \overline{\overline{Zy9}}$$

$$sYa1 = Iyp C10 + \overline{E_y} \overline{Ya1} \underline{Ya2 \overline{Ya3 \overline{Ya4 \overline{E_y}}}}$$

$$rYa1 = Iyc + \overline{E_y} \overline{Ya1} \underline{Ya2 \overline{Ya3 \overline{Ya4 \overline{E_y}}}}$$

$$Zy10 = \overline{\overline{Zy10}}$$

$$Zy11 = \overline{\overline{Zy11}}$$

$$sYa0 = Iye C18 + \overline{E_y} \overline{Ya0} \underline{Ya1 \overline{E_y}}$$

$$rYa0 = Iyc$$

$$Zy12 = \overline{\overline{Zy12}}$$

$$Zy13 = \overline{\overline{Zy13}}$$

3.270 Pin Address Counter

$$sYpa = Ioc C17 C13 \overline{C19} \overline{C20} \overline{C21} \overline{C22} \overline{C23} T0$$

$$rYpa = Pin T0 + St$$

$$Zy14 = \overline{\overline{Zy14}}$$

$$Zy15 = \overline{\overline{Zy15}}$$

3.271 Skip Gate

$$Skr = \overline{C1} C17 \overline{C9} \overline{C10} C11 \overline{C19} \overline{C20} \overline{C21} \overline{C22} \overline{C23} C15 Ysc^*$$

$$+ \overline{C1} C17 \overline{C9} \overline{C10} C11 \overline{C19} \overline{C20} \overline{C21} \overline{C22} \overline{C23} C12 \overline{Y10} \overline{Y11} \overline{Y12} \overline{Y13} \overline{Y14}^*$$

$$+ \overline{C1} C17 \overline{C9} \overline{C10} C11 \overline{C19} \overline{C20} \overline{C21} \overline{C22} \overline{C23} C13 Iyf^*$$

$$+ \overline{C1} C17 \overline{C9} \overline{C10} C11 \overline{C19} \overline{C20} \overline{C21} \overline{C22} \overline{C23} C14 \overline{Y_e}^*$$

$$+ C10 \overline{C11} C13 \overline{Y10} \overline{Y11} \overline{Y12} \overline{Y13} \overline{Y14} \overline{C1}^*$$

$$Zy16 = \overline{\overline{Zy16}}$$

$$Zy17 = \overline{\overline{Zy17}}$$

$$Zy18 = \overline{\overline{Zy18}}$$

$$Zy19 = \overline{\overline{Zy19}}$$

$$Zy20 = \overline{\overline{Zy20}}$$

$$Zy21 = \overline{\overline{Zy21}}$$

$$Zy22 = \overline{\overline{Zy22}}$$

$$Zy23 = \overline{\overline{Zy23}}$$

Yhs	=	$\overline{\text{Yhs}}$		Ry14	=	Ry14
Yes	=	$\overline{\text{Yes}}$		Ry15	=	Ry15
$\overline{\text{Mtg}}$	=	$\overline{\text{Mtg}}$		Ry16	=	Ry16
$\overline{\text{Npy}}$	=	$\overline{\text{Npy}}$		Ry17	=	Ry17
Ecy	=	$\overline{\text{Ecy}}$		Ry18	=	Ry18
Yx12	=	$\overline{\text{Yx12}}$	(93221 only)	Ry19	=	Ry19
Yx24	=	$\overline{\text{Yx24}}$	(93221 only)	Ry20	=	Ry20

3. 273 Input/Output Signals Generated

Ryp	=	Ryp		Ry22	=	Ry22
Ry1	=	Ry1		Ry23	=	Ry23
Ry2	=	Ry2		Ry24	=	Ry24
Ry3	=	Ry3		Y9	=	Y9
Ry4	=	Ry4		Y10	=	Y10
Ry5	=	Ry5		Y11	=	Y11
Ry6	=	Ry6		Y12	=	Y12
Ry7	=	Ry7		Y13	=	Y13
Ry8	=	Ry8		Y14	=	Y14
Ry9	=	Ry9		Y0	=	Y0
Ry10	=	Ry10		Y5	=	Y5
Ry11	=	Ry11		Y6	=	Y6
Ry12	=	Ry12		Iy	=	Iy
Ry13	=	Ry13				

SECTION IV
INSTALLATION AND MAINTENANCE

4. 1 GENERAL

4. 2 This section contains information relating to the installation and maintenance of Model 932XX series TMCCs. As the Model 922XX series TMCCs are no longer being installed, only the Model 932XX series is covered in this section.

4. 3 INSTALLATION

4. 4 The basic 925/930/9300 computers are shipped with the TMCC physically installed. After installation of the computer, the intercabling of the TMCCs must be performed.

4. 5 INTERCABLING

4. 6 Figure 4-1 illustrates typical intercabling of the Model 93200 TMCC for the 925/930/9300 computers. Figure 4-2 illustrates the intercabling of the Model 93221 TMCC. Power distribution for the various chassis is illustrated in figure 4-3.

4. 7 Intercabling of the input/output devices to the W (or A) channel may be found in the applicable input/output device technical manual.

4. 8 After intercabling the TMCCs and the input/output devices, a program should be run to ensure proper operation of the W (or A) channel. Any diagnostic program utilizing the input/output device may be run.

4. 9 925/930 COMPUTER W CHANNEL TEST PROGRAM

4. 10 Table 4-1 lists a sample program which may be run to check out the W channel for proper operation. This test program causes the message ASSEMBLY DONE ENTER NEW PROGRAM to be typed out under program control. The computer stores the internal codes for these characters in memory beginning in location 2000. The routine inserts the carriage return code, 52, and the space code, 12, where needed and requests End-of-Record interrupt. It is written as a closed subroutine using interrupts, channel W and Typewriter Number One. The internal code for the output message is as follows:

A	S	S	E	M	B	L	Y	Sp	D	O	N	E	C/R	E	N	2000
21	62	62	25	44	22	43	71	12	24	46	45	25	52	25	45	
T	E	R	Sp	N	E	W	Sp	P	R	O	G	R	A	M	Sp	2004
62	25	51	12	45	25	66	12	47	51	46	27	51	21	44	12	

4. 11 9300 COMPUTER A CHANNEL TEST PROGRAM

4. 12 Table 4-2 lists a sample program which may be run to check out the A channel for proper operation. This test program causes the message ASSEMBLY DONE ENTER NEW PROGRAM to be typed out under program control. The computer internal codes for these characters are stored beginning in location 02000. The carriage return code, 052, and the space code, 012, are inserted where needed. The End-of-Record interrupt is requested. The routine is written as a closed subroutine which uses interrupts, channel A, and Typewriter Number One. The internal code for the output message is the same as given in paragraph 4. 10.

4. 13 MODULE LOCATION

4. 14 Figure 4-4 illustrates the location of all modules for the various models of TMCCs.

4. 15 MAINTENANCE

4. 16 The following information is presented as an aid in maintaining the Models 932XX TMCCs. Presented herein are descriptions and timing diagrams of the signals available on the various input/output connectors and the diagnostic test programs for maintenance of the TMCC.

4. 17 PERIODIC INSPECTION

4. 18 No periodic inspection is required for the TMCC other than that required for the computer as a whole. No attempt should be made to periodically check for loose wires, poor solder connections, or bent pins because of the packaging density of the wiring and components and the possibility of causing malfunctions. Wiring layout and length is critical in some areas and should not be touched except for correcting a malfunction.

4. 19 CORRECTIVE MAINTENANCE

4. 20 If it should become necessary to replace a component on a module, the physical location of the component, quantity, type, and part number are indicated on the module drawings contained in Section 6.

4. 21 INPUT/OUTPUT SIGNALS AND TIMING RELATIONSHIPS

4. 22 The signals described are available on the input/output connectors as illustrated in figure 4-5. The signals are theoretical and do not take into consideration circuit and transmission delays which tend to add 200 to 400 nanoseconds of delay.

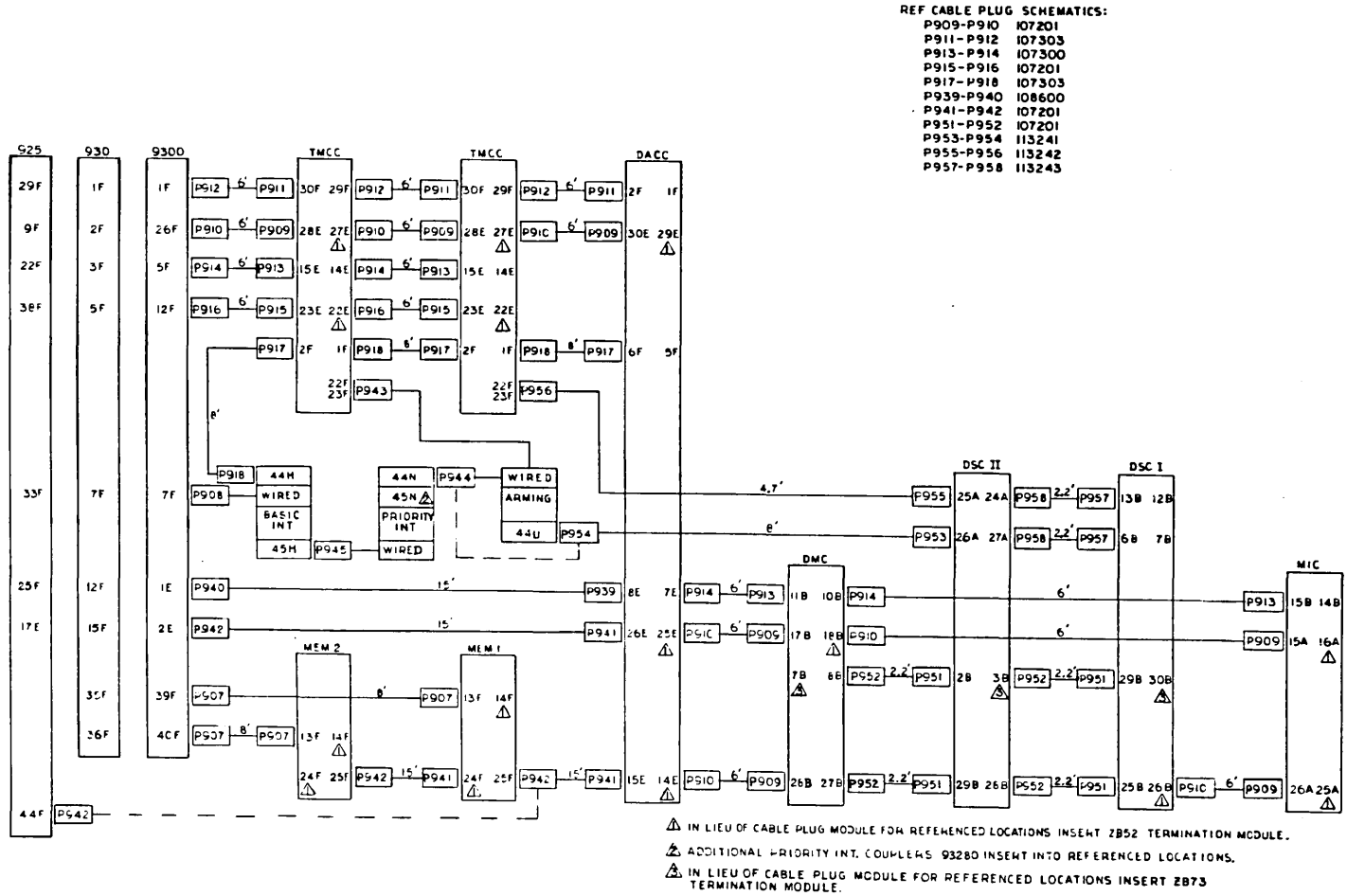


Figure 4-1. Model 93200 TMCC, Intercabling Diagram

REF CABLE PLUG SCHEMATICS:
 P909-P910 I07201
 P911-P912 I07303
 P913-P914 I07300
 P915-P916 I07201
 P917-P918 I07303
 P939-P940 I08600
 P941-P942 I07201
 P951-P952 I07201
 P953-P954 I13241
 P955-P956 I13242
 P957-P958 I13243

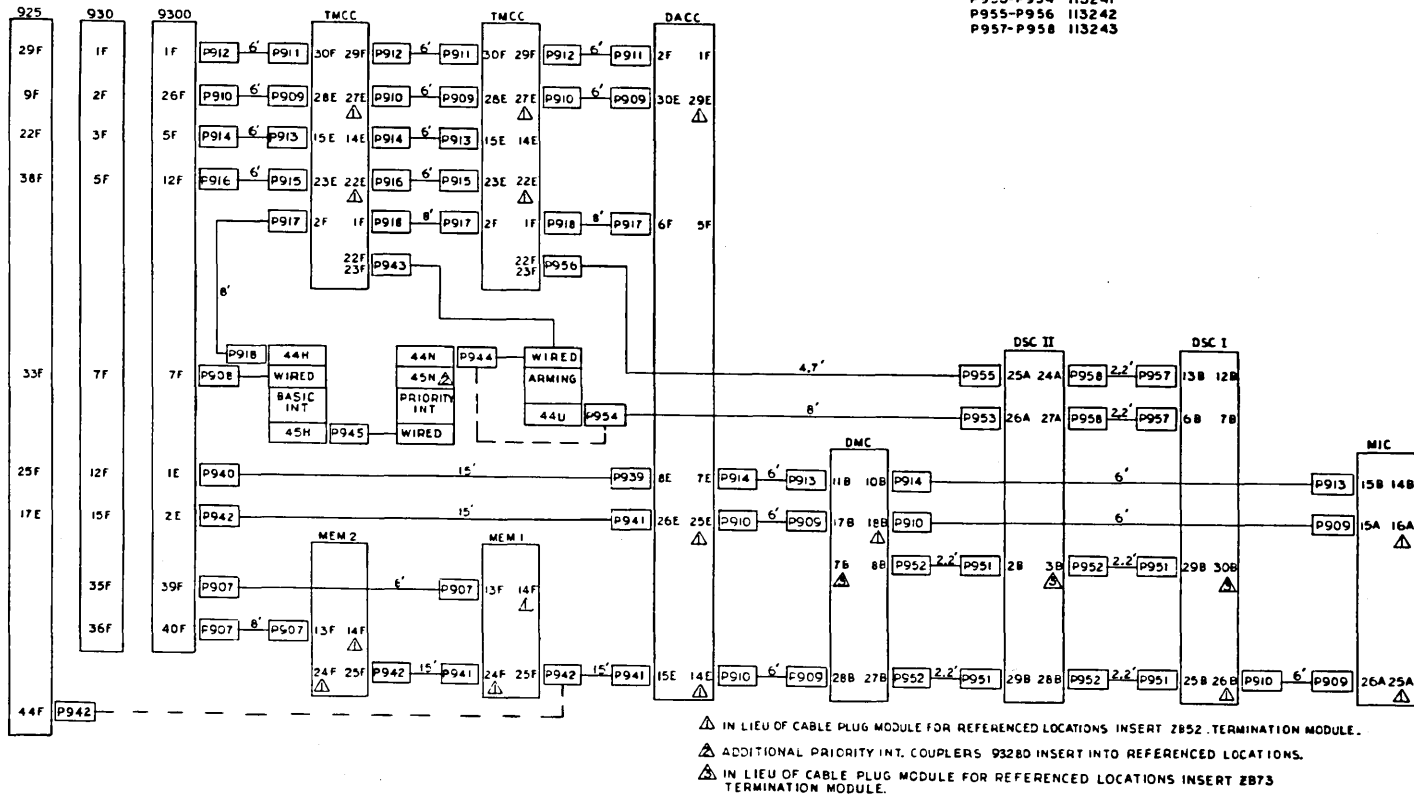


Figure 4-2. Model 93221 TMCC, Intercabling Diagram

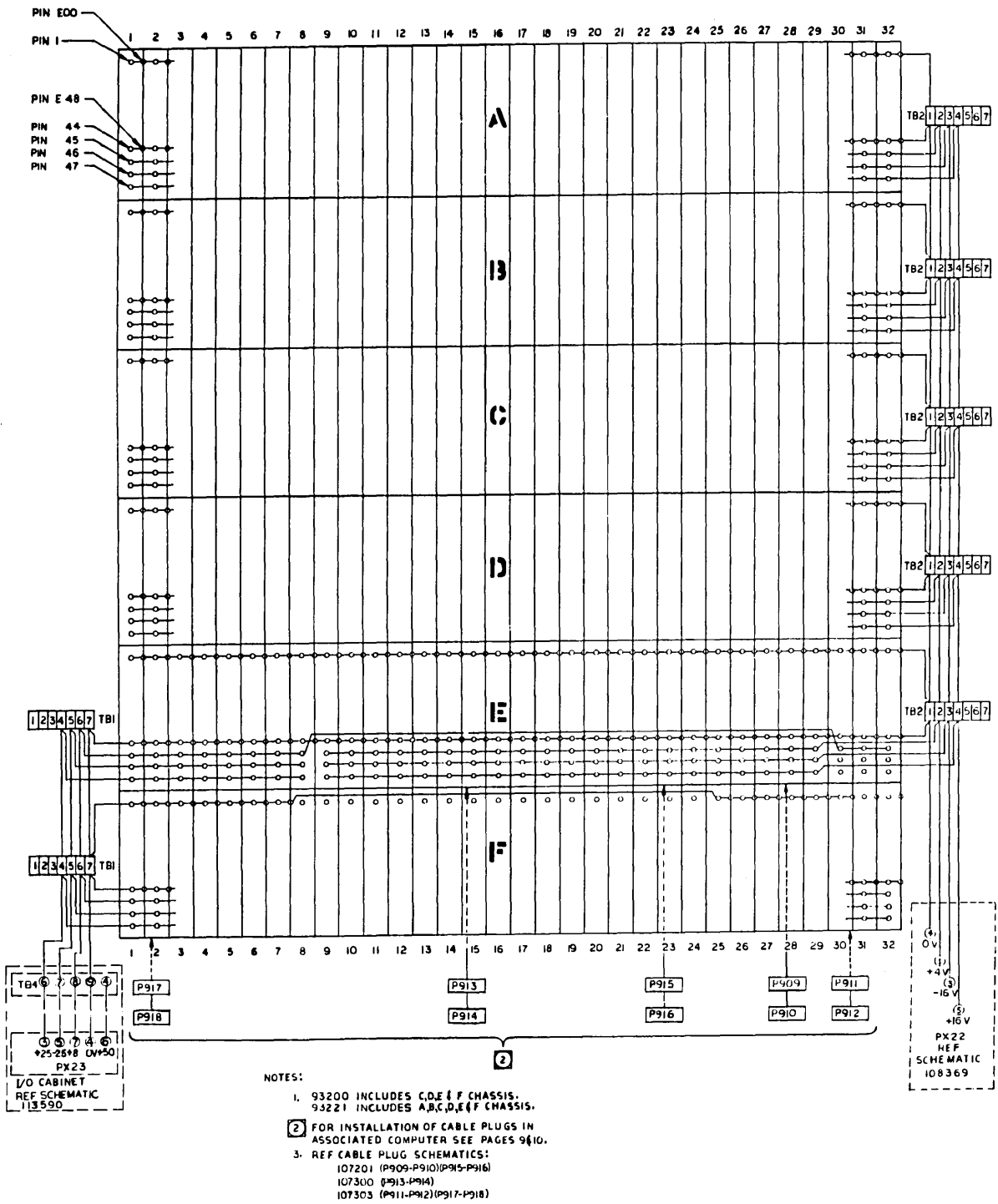


Figure 4-3. Power Distribution Diagram

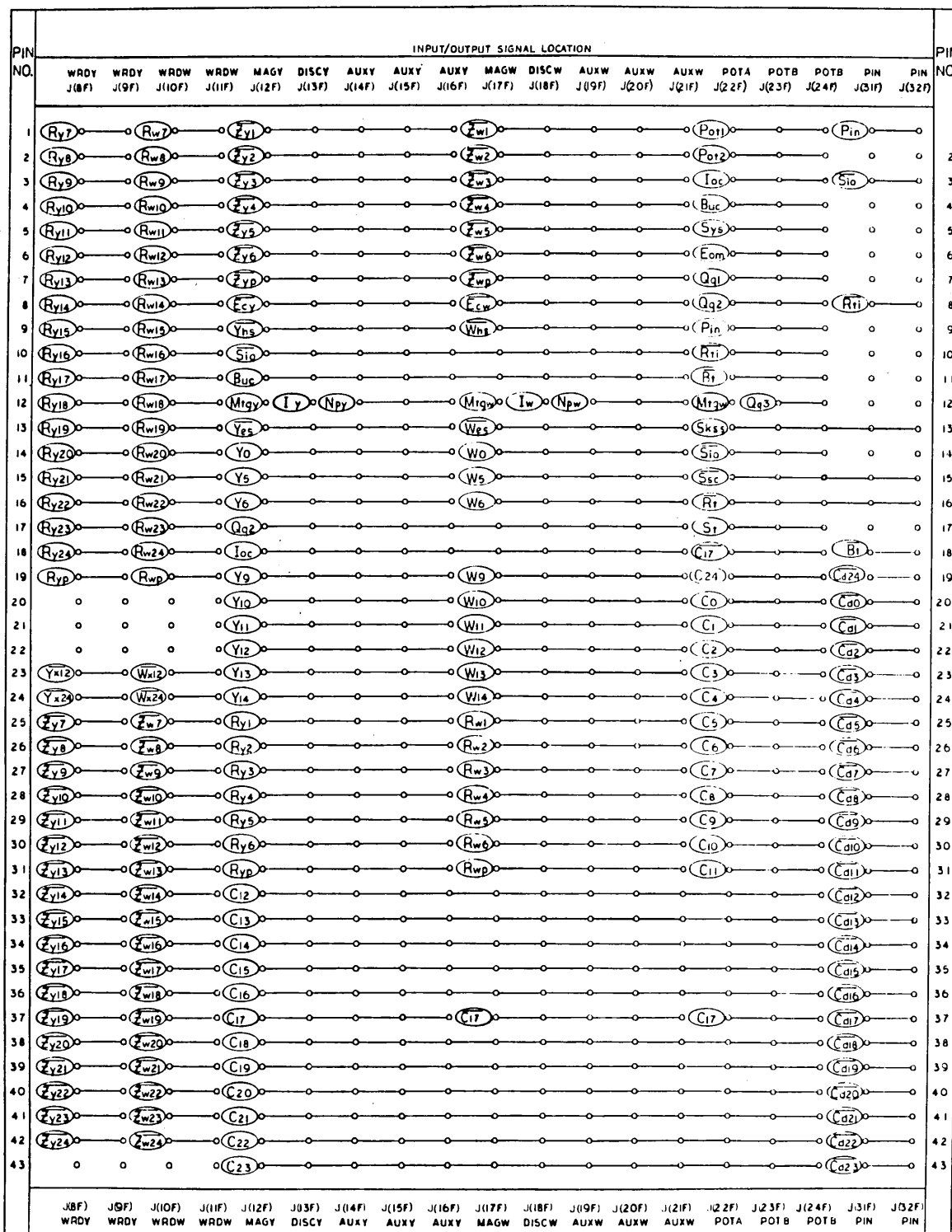


Figure 4-5. Input/Output Signal Location Diagram

Table 4-1. 925/930 Computers, W Channel Sample Test Program

Location	Instruction	Address	Comments
1000	PZE		This instruction is an assembler instruction, used as a convenient way to reserve the entry location for subroutine use.
	CLR		Clears the A and B Registers.
	STA	SWICH	Clears the location called SWICH. SWICH later indicates to the main program that output is complete.
	TYP	*0, 1, 4	Connects Typewriter Number One to channel W for output, specifies four characters per word mode, and alerts channel W interlace. The instruction is an EOM with octal configuration, 0 02 40641.
	EXU	WRITE	Causes the Input/Output EOM in location WRITE to be executed.
	POT	WRITE + 1	Sends the word count and starting address in WRITE + 1 to the channel.
	BRR	1000	Branches back to the main program.
WRITE	EOM 00403720	16200	Specifies output function code 00 and the End-of-Record interrupt. The word in WRITE + 1 specifies that eight words will output from memory beginning in location 2000. According to output function 00, when the word count equals zero during the transmission, the output terminates, and when the last character is out, the device disconnects and the interrupt occurs.
33	BRM	OKAY	Branches and marks to location OKAY elsewhere in memory.
OKAY	PZE		Saves the entry location
	MIN	SWICH	Increments location SWICH as an indicator for the main program.
	BRU	*OKAY	Branches to the main program and clears the active interrupt.

Table 4-2. 9300 Computer, A Channel Sample Test Program

Location	Instruction	Address	Comments
01000	PZE		This instruction is an assembler instruction used as a convenient way to reserve the entry location for subroutine use.
	STZ	SWICH	Clears the location called SWICH. SWICH is later used to indicate to the main program that output is complete.
	TYP	*0, 1, 4	Connects Typewriter Number One to channel A for output, specifies four characters per word mode, and alerts channel A interlace. The instruction is an EOM with octal configuration, 0 02 42641.
	EXU	WRITE	Causes the Input/Output Control EOM in location WRITE to be executed.
	POT	WRITE + 1	Sends the word count and starting address in WRITE + 1 to the channel.
	BRR	01000	Branches back to the main program.
WRITE	EOM 00403720	016200	Specifies output function code 01 (IOSD) and the End-of-Record interrupt. The word in WRITE + 1 specifies that eight words will be output from memory beginning in location 03720. According to output function 01 (IOSD) when the word count equals zero during the transmission, the device is disconnected when the last character is out and the interrupt then occurs.
011	BRM	OKAY	Branches and marks to location OKAY elsewhere in memory.
OKAY	PZE		Saves the entry location.
	MPO	SWICH	Increments location SWICH as an indicator for the main program.
	BRC	*OKAY	Branches to the main program and clears the active interrupt, level 011.

4.23 The logic terms for signals generated in the main frame referred to in these paragraphs represent the 930 computer logic equations. Although the 925 and 9300 computer implementation may differ somewhat from the 930 computer, the functions achieved are similar.

4.24 $\overline{Qq1}$, $\overline{Qq2}$, and $\overline{Qq3}$ (See figure 4-6)

4.25 $\overline{Qq1}$, $\overline{Qq2}$, and $\overline{Qq3}$ are clocking signals provided to external devices. They are functionally similar to the Q1 and Q2 signals provided to external equipment on the 910 and 920 computers. Signals $\overline{Qq1}$, $\overline{Qq2}$, and $\overline{Qq3}$ are derived from the Pulse Counter, Qr1 through Qr4, in the TMCC.

$$\overline{Qq1} = T_5 - T_0$$

$$T_5 - T_0 = Q_{r3} \overline{Q_{r1}} \overline{Q_{r4}}$$

$$\overline{Qq2} = T_6 - T_3$$

$$T_6 - T_3 = Q_{r3} Q_{r4} + Q_{r3} \overline{Q_{r2}} \overline{Q_{r1}}$$

$$\overline{Qq3} = T_7 - T_4$$

$$T_7 - T_4 = Q_{r4}$$

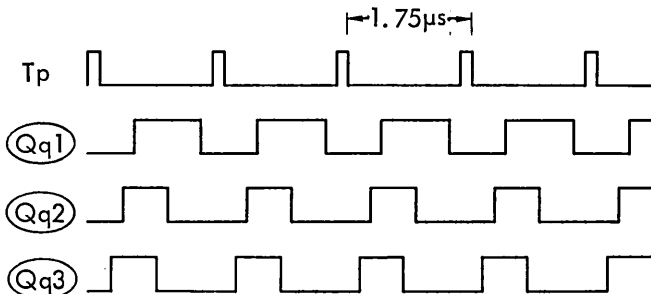


Figure 4-6. 930/9300 Timing Diagram, $\overline{Qq1}$, $\overline{Qq2}$, $\overline{Qq3}$

4.26 \overline{Eom} , \overline{Buc} , \overline{Ioc} , \overline{Sys} (See figure 4-7)

4.27 \overline{Eom} is the execution signal for the EOM instruction.

$$\overline{Eom} = Eom \overline{Er}$$

The term \overline{Er} inhibits \overline{Eom} during the interlace loading sequence. \overline{Eom} is derived from the CPU and transmitted to the TMCC. \overline{Eom} is true from T7 through Tr.

$$\overline{Eom} = Q_5 \overline{01} \overline{04} \overline{05} \overline{T_s} (Q_2 + Q_5) + \dots$$

$$Q_2 + Q_5 = T_7 - Tr$$

During FILL operations, a pseudo Eom is generated.

$$\overline{Eom} = I_x \overline{Go} \overline{Ht} \overline{Kg} (A_2 + Q_5) + \dots$$

The term $\overline{T_s}$ inhibits \overline{Eom} during all time share operations.

4.28 \overline{Buc} is a control signal derived from the EOM instruction and is used to activate the TMCC and peripheral devices. \overline{Buc} is true from T7 through Tr.

$$\overline{Buc} = Eom \overline{C10} \overline{C11} \overline{C1}$$

For TMCC channels C and D,

$$\overline{Buc} = Eom \overline{C10} \overline{C11} C1$$

Peripheral devices must use $\overline{C17}$ and C17 to distinguish between channels W and Y and similarly, between channels C and D.

4.29 \overline{Ioc} is an input/output control signal derived from the EOM instruction. \overline{Ioc} is true from T6 through T0.

$$\overline{Ioc} = Eom \overline{C10} C11 \overline{C1} \overline{Er} Q_{r3}$$

$$Q_{r3} = T_6 - T_0$$

For channels C and D,

$$\overline{Ioc} = Eom \overline{C10} C11 C1 \overline{Er} Q_{r3}$$

The term \overline{Er} is used to inhibit \overline{Ioc} during the interlace loading sequence. Peripheral devices must use $\overline{C17}$ and C17 to distinguish between channels W and Y, or similarly, between channels C and D.

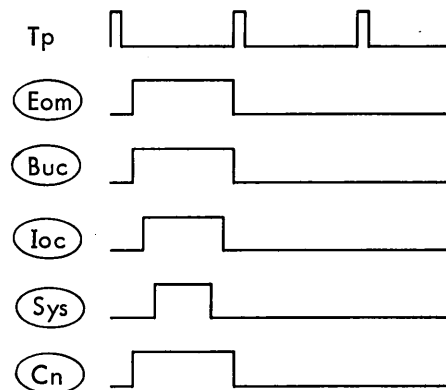


Figure 4-7. 930/9300 Timing Diagram, \overline{Eom} , \overline{Buc} , \overline{Ioc} , \overline{Sys}

4.30 Sys is a control signal for systems communication derived from an EOM instruction. Sys is true from T5 through T1.

$$\text{Sys} = \text{Eom C10 C11 } \overline{\text{C9}} \overline{\text{Tsr}} (\text{T5} - \text{T1})$$

Systems devices must use $\overline{\text{C17}}$ and C17 to distinguish between channels W and Y, or similarly, between channels C and D. The term $\overline{\text{Tsr}}$ inhibits Sys during time-share operations with the TMCC.

4.31 Pin , $\overline{\text{Rt}}$, $\overline{\text{Cdn}}$, $\overline{\text{Rti}}$
(See figures 4-8, 4-9, and 4-10)

4.32 The PIN instruction permits direct parallel entry of up to 24-bits of data to memory via the C register. A "Ready for Input" signal, Pin , is provided by the TMCC to external equipment.

$$\begin{aligned} \text{Pin} &= \text{Pin} \\ \text{Pin} &= \overline{\text{F1}} \overline{\text{F2}} \overline{\text{F3}} \text{02} \text{06} \overline{\text{T5}} \text{Q1} \\ \text{Q1} &= \text{T7} - \text{T0} \end{aligned}$$

While the Pin signal is true, the C register is first reset and then the external data, $\overline{\text{Cdn}}$, is strobed.

$$\begin{aligned} \text{rCn} &= \text{Cxi Q2} \\ \text{Cxi} &= \overline{\text{F1}} \overline{\text{F2}} \overline{\text{F3}} \text{02} \text{06} \overline{\text{T5}} \text{Q1} \\ \text{Q2} &= \text{T7} - \text{T3} \\ \text{sCn} &= \text{Cxi } \overline{\text{Cdn}} \\ \overline{\text{Cdn}} &= \overline{\overline{\text{Cdn}}} \end{aligned}$$

The process of resetting the C register and then strobing the data repeats until the external device provides a "Ready" indication by making $\overline{\text{Rt}}$ false. The data lines, $\overline{\text{Cdn}}$, must be in a stable condition prior to the time $\overline{\text{Rt}}$ goes false and should remain stable for the duration of the Pin signal. Until $\overline{\text{Rt}}$ goes false, the PIN instruction was locked in phase Ø2 , but is now permitted to proceed to phase Ø4 . A "PIN Complete" signal, $\overline{\text{Rti}}$, is generated and sent to the external device.

$$\begin{aligned} \text{sRf} &= \text{06 F2 } \overline{\text{F3}} \text{Q2 } \overline{\text{Rt}} + \dots \\ \overline{\text{Rt}} &= \overline{\overline{\text{Rt}}} \\ \text{sF1} &= \overline{\text{F1}} \overline{\text{F3}} \text{01} \text{03} \text{04} \overline{\text{Ia}} \overline{\text{Rf}} \text{Tp} + \dots \end{aligned}$$

$$\begin{aligned} \text{rF2} &= \overline{\text{F1}} \overline{\text{F3}} \text{01} \text{03} \text{04} \overline{\text{Ia}} \overline{\text{Rf}} \text{Tp} + \dots \\ \overline{\text{Rti}} &= \overline{\text{Ø4} \text{01} \text{04} \text{06} \overline{\text{T5}} (\text{Q2} + \text{Q5})} \end{aligned}$$

$$(\text{Q2} + \text{Q5}) = \text{T7} - \text{Tr}$$

$$\overline{\text{Rti}} = \overline{\overline{\text{Rti}}}$$

If $\overline{\text{Rt}}$ is always held false during the PIN instruction, the PIN instruction remains in phase Ø2 for only one cycle.

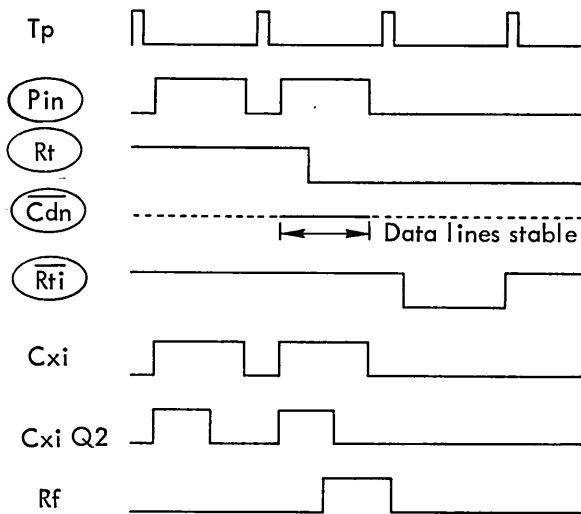


Figure 4-8. 930/9300 Timing Diagram, Pin , $\overline{\text{Rt}}$, $\overline{\text{Cdn}}$, $\overline{\text{Rti}}$

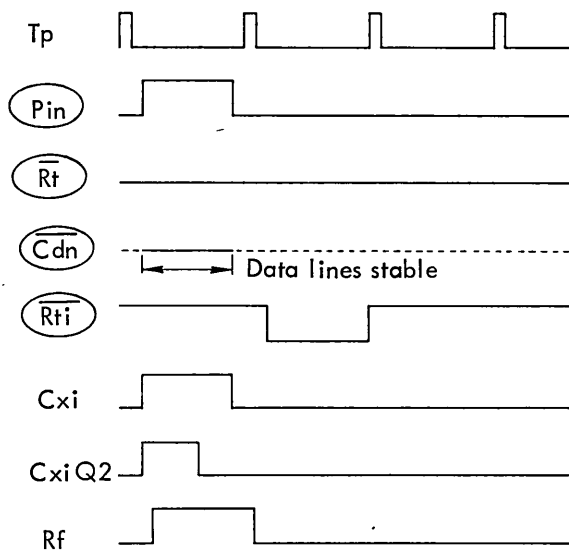


Figure 4-9. 930/9300 Timing Diagram, Pin , $\overline{\text{Rt}}$, $\overline{\text{Cdn}}$, $\overline{\text{Rti}}$, $\overline{\text{Rt}}$ Initially False

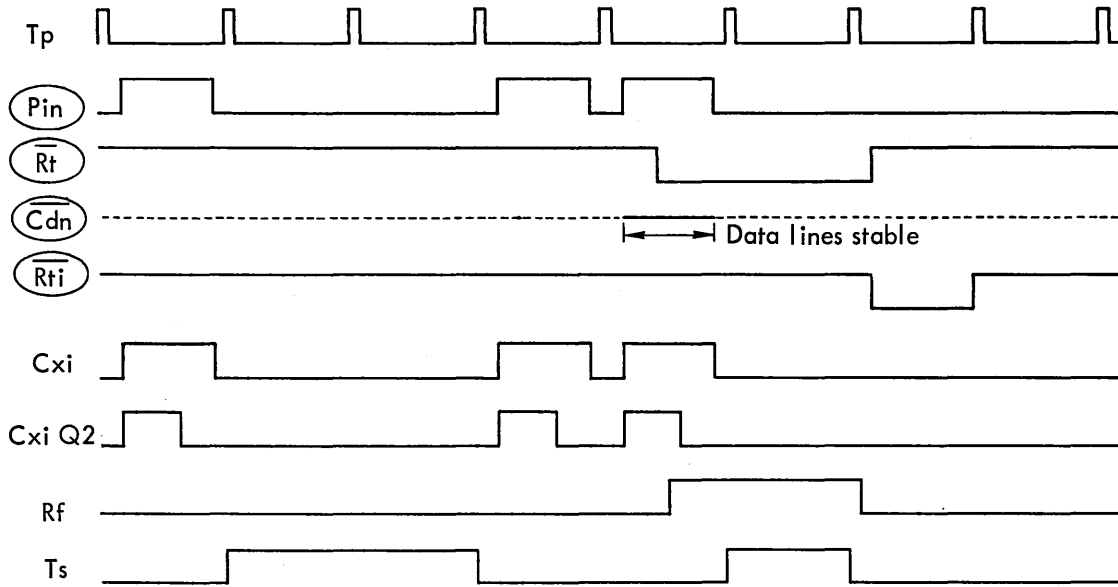


Figure 4-10. 930/9300 Timing Diagram, \overline{Pin} , \overline{Rt} , \overline{Cdn} , \overline{Rti} , Effects of Time-Share

4.33 Should a time-share operation occur during phase $\emptyset 2$ of the PIN instruction, the \overline{Pin} signal is inhibited. At the completion of the time-share operation, depending on the condition of \overline{Rt} , at least one more cycle of the \overline{Pin} signal occurs. Should a time-share operation occur during phase $\emptyset 4$, the "Pin Complete" signal is inhibited until the time-share operation(s) is/are completed.

4.34 $\overline{Pot 1}$, $\overline{Pot 2}$, \overline{Rt} , \overline{Cn}
 (See figures 4-11, 4-12, and 4-13)

4.35 The POT instruction permits direct parallel output of up to 24-bits of data to external devices from memory via the C register. The $\overline{Pot 1}$ signal denotes to external devices that the computer is in the process of executing a POT instruction.

$$\overline{Pot 1} = \overline{Pot 1}$$

$$Pot 1 = \overline{F1} \overline{F2} \overline{F3} \overline{O2} \overline{O6}$$

The $\overline{Pot 2}$ signal denotes to external devices that the contents of the C register may be strobed.

$$\overline{Pot 2} = \overline{Pot 1} (T5 - T1) \overline{Tsr}$$

The $\overline{Pot 1}$ signal stays true and the $\overline{Pot 2}$ signal repeats as long as the computer is locked in phase $\emptyset 2$.

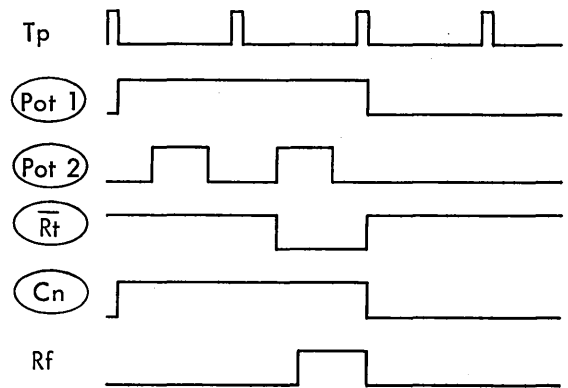


Figure 4-11. 930/9300 Timing Diagram, $\overline{Pot 1}$, $\overline{Pot 2}$, \overline{Rt} , \overline{Cn}

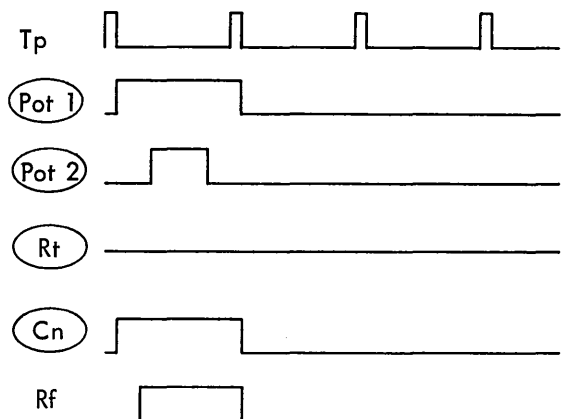


Figure 4-12. 930/9300 Timing Diagram, $\overline{Pot 1}$, $\overline{Pot 2}$, \overline{Rt} , \overline{Cn} , \overline{Rt} Initially False

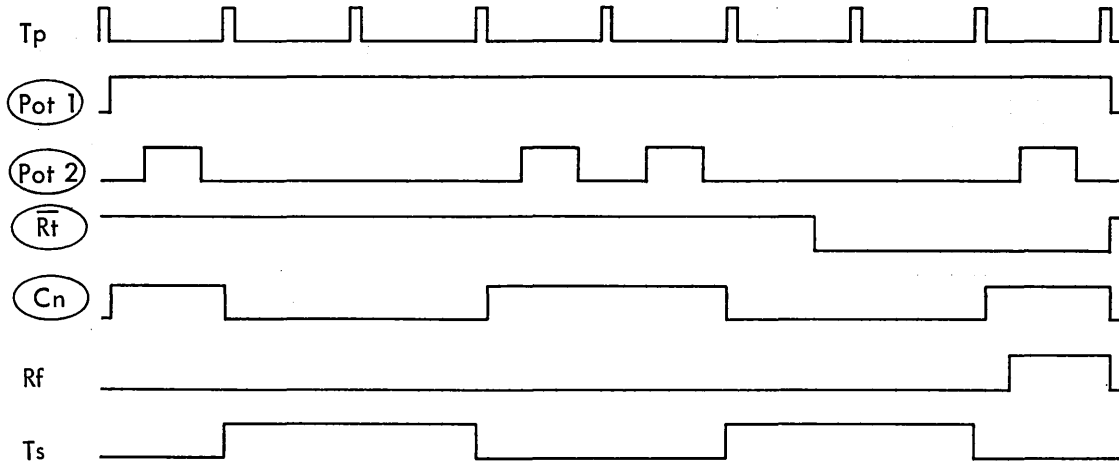


Figure 4-13. 930/9300 Timing Diagram, (Pot 1), (Pot 2), (Rt), (Cn) Effects of Time-Share

When the Ready signal, (Rt), from the external device goes false (low), the POT instruction proceeds to phase 06. (Pot 1) and (Pot 2) are then inhibited.

$$sRf = 06 F2 \overline{F3} Q2 Rt$$

$$Rt = (Rt)$$

$$sF1 = \overline{F1} \overline{F3} \overline{01} \overline{03} \overline{04} \overline{Ia} Tp Rf$$

If (Rt) is always held false during the POT instruction, the POT instruction remains in phase 02 for only one cycle.

4.36 Should a time-share operation occur during the phase 02, the (Pot 1) signal remains true but (Pot 2) is inhibited. At the completion of the time-share operation, depending on the condition of (Rt), at least one more cycle of (Pot 1) and (Pot 2) occurs. It is necessary that (Rt) be held at ground until the computer acknowledges receipt of the Ready signal by making (Pot 1) go false.

4.37 (Skss), (Ssc) (See figure 4-14)

4.38 On instructions for external system devices, an Sks strobe pulse, (Skss), is provided from the TMCC.

$$(Skss) = Skss$$

$$Skss = 05 01 \overline{04} \overline{Ts} (\overline{A00} (Q3 + Q5) + \overline{C9})$$

$$Q3 + Q5 = T6 - Tr$$

The SKS instruction remains in phase 05 for two cycles to permit the C register outputs, (Cn), to attain a

stable configuration at the external system device. The signal (Skss) is generated during the second cycle if C9 is true. (Skss) is true for approximately two cycles if C9 is true. The response from external system devices, (Ssc), (false for a skip condition) is received by the TMCC and is sent to the computer as Skrz. If a skip is to occur, the Sk flip-flop is set at Tr time during the second cycle of phase 05.

$$sSk = 05 01 \overline{04} \overline{A00} Tr Sks + \dots$$

$$Sks = Skrz + \dots$$

If a time-share operation occurs during the SKS instruction, phase 05 is repeated for two cycles, thereby permitting the C register outputs, (Cn), to attain a stable configuration prior to the generation of the (Skss) probe pulse.

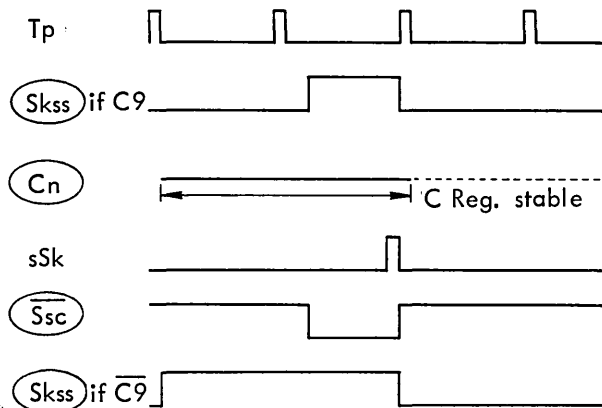


Figure 4-14. 930/9300 Timing Diagram, (Skss), (Cn), (Ssc)

4.39 If $\overline{C9}$ is true, the Skss signal may occur for only one cycle if a time-share operation occurs during the SKS instruction. At the completion of the time-share, however, a two cycle Skss signal occurs. If $C9$ is true and time-share operation occurs, the Skss signal occurs only once, during the last of the two phase $\phi 5$ cycles.

4.40 TEST PROGRAMS

4.41 Proper operation of the TMCC can be checked by performance of the applicable test program. Operation in the compatible mode of the TMCC can be checked by performing the test program for the input/output device connected to the TMCC. The sample test programs given in tables 4-1 and 4-2 may be performed. Test programs for the extended modes are given in the following paragraphs.

4.42 EXTENDED MODE I/O TEST PROGRAM FOR 925/930 COMPUTERS

4.43 This program tests as many of the extended I/O operations as possible with paper tape. Any 925/930 computer with a typewriter attached to the W channel and a paper-tape punch and reader on any interlaced communication channel may be utilized. The W channel need not be interlaced for the typewriter.

4.44 The program occupies 838₁₀ locations from 0177 to 1714 octal. It is supplied on paper tape with a self-loading bootstrap. To load, insert tape in tape reader and perform fill procedure.

4.45 Fill

4.46 The fill procedure is as follows:

- a. Set up selected input device with the input program. Initial portion of the program contains the "bootstrap" program.
- b. Set RUN-IDLE-STEP switch to the IDLE position.
- c. Press START switch.
- d. Press PAPER TAPE FILL switch. This causes a WIM 2 (03200002) instruction to be inserted into the instruction register and loads the Index Register with 77777771. The FILL switch also prepares the channel to operate in the forward, binary, four characters per word mode.

4.47 Operation

4.48 To select the reader and punch units to be used enter:

- "CU" P to select punch
- "CU" R to select reader

The letter C is the channel number and may be any digit 0-7, and the letter U is the unit, either one or two. The reader and punch need not be on the same channel.

4.49 The test is started by entering the letter "S". Control of the test operation is then a function of the Breakpoint Switches. Table 4-3 summarizes the switch functions.

Table 4-3. 925/930 Computers Breakpoint Switch Functions

Breakpoint Switch	Reset	Set
1	Run in the normal mode as determined by the other switches	Stop and return to keyboard control at the end of the current pass (punch or read)
2	Continue to run test selected by switch 3	Cycle test runs from punch to read to punch and so on
3*	Selects punch mode	Selects read mode
4	Stop and type diagnostic messages whenever an error occurs	Do not stop and type on errors but continue to run.

*Used when switch 2 is reset or when starting test

4.50 When running cyclic tests from punch to read, the tape from the punch should be inserted into the reader.

4.51 Punching

4.52 The program punches four blocks of 64 characters each in one pass. The characters form a counting sequence from 00 to 77g. The first block is started with leader and output with an IOSP. All punching is done in the one character per word mode. When the word count reaches zero, an IOSD is loaded to punch a second block of 64 characters. No leader is punched between the first and second blocks. This results in one physical block 128 characters long. Starting with leader, two additional blocks of 64 characters are then punched with an IOSD.

4.53 At the conclusion of each output operation, the channel address register is stored and compared with the expected value. If they do not agree, the program types the expected and actual values.

4.54 The program tests the channel during the output operation to see if the channel should erroneously disconnect before the word count reaches zero.

4.55 Reading

4.56 Each of the four blocks is read with a different set of commands and counts so as to test as many operations as possible. After reading a block, a general subroutine checks for input parity errors, channel end address for agreement with the expected address, and the data read character by character. Error messages with block numbers are typed in the event of any one of these tests failing. If a test fails, reference should then be made to the test program flow diagram and troubleshooting information in Section 5 of this manual. The handling of each block is as follows:

a. Block 1. The first block is one-half of a 128 character physical block. Reading one character per word, an IOSD with a count of 64 is used to read this block. The program checks to see if the count reaches zero and the channel becomes inactive at the same time.

b. Block 2. This is the second half of the first physical block and is read with an IOSP with a count of 65. The read should terminate because of the end of record. The program checks to see that the word count does not reach zero and the channel remains active after the CIT (inter-record test) instruction skips. The tape is finally stopped with a disconnect before the data is checked.

c. Block 3. Block 3 is a 64-character physical block. It is read with two channel commands. The

first is an IOSP with a count of 32. If the count goes to zero before the channel disconnects, an IORP with a count of 33 is loaded. This should cause the inter-record indicator to be turned on at the end of the record. The count should not reach zero and the channel should remain active. The tape is again stopped with a disconnect before the data is checked.

d. Block 4. This is the third physical block of 64-characters and is read with an IORD with a count of 56. The program waits for the channel to be inactive then checks to see if the channel ignored the last eight characters. If the tape was erroneously stopped after the 56th character it will show up as a failure on the first block of the next read pass.

4.57 Test Program

4.58 Table 4-4 gives the test program for the 925/930 computers I/O extended mode.

4.59 EXTENDED MODE I/O TEST PROGRAM FOR 9300 COMPUTER

4.60 This program tests as many of the extended I/O operations as possible with paper tape. Any 9300 computer with a typewriter attached to the W channel and a paper-tape punch and reader on any interlaced communication channel may be utilized. The W channel need not be interlaced for the typewriter.

4.61 The program occupies 838₁₀ locations from 0177 to 1714g. It is supplied on paper tape with a self-loading bootstrap. To load, insert tape in tape reader and perform fill procedure.

4.62 Fill

4.63 The fill procedure is as follows:

a. Press POWER switch on. When power is on, the switch is lighted.

b. Press IDLE switch.

c. With computer in IDLE, press RESET switch. This clears the D register and the program counter.

d. Press RUN switch. The computer now executes the instruction in the D register (which is a HALT instruction). The program counter advances to 00001.

e. Press PAPER TAPE LOAD switch. This switch causes an AIM 2 (0 32 00002) instruction to be inserted into the D register and clears the HALT instruction. Index 1 is loaded with 001 77771. The LOAD switch also prepares the channel to operate in the forward, binary, four character per word mode.

Table 4-4. 925/930 Computers, Extended Mode I/O Test Program (Sheet 1 of 20)

80001642

```

1 TYPE EQU 930 MACHINE DEFINITION: 930/925
2 *
3 *
4 * EXTENDED MODE I/O TEST PROGRAM. (925/930/9300 ALL CHANNELS)
5 *
6 * A.W. ENGLAND, SDS
7 *
8 *
9 * THIS PROGRAM USES THE PAPER TAPE READER AND PUNCH TO CHECK THE
10 * OPERATION OF THE I/O CHANNELS IN THE EXTENDED INTERLACE MODE OF
11 * OPERATION. THE READER AND PUNCH MAY BE CONNECTED TO ANY INTERLACED
12 * TMCC OR DACC. THEY NEED NOT BE ATTACHED TO THE SAME CHANNEL. THE
13 * PROGRAM ADDRESSES THE KEYBOARD FOR INFORMATION ABOUT UNIT AND CHANNEL
14 * SELECTION.
15 *
16 * TO SELECT PUNCH TYPE: *CU*P
17 *
18 * TO SELECT READER TYPE: *CU*R
19 *
20 * WHERE *C* REPRESENTS CHANNEL NO. AND MAY HAVE THE VALUES 0-7;
21 * AND *U* REPRESENTS UNIT NO. AND MAY HAVE THE VALUE 1-2. THE LETTER
22 * P OR R CAUSE THE SELECTION TO BE MADE WITH THE TWO PREVIOUS DIGITS.
23 * SPACES SHOULD NOT BE TYPED BETWEEN DIGITS OR CONTROL CHARACTER.
24 *
25 * TO START THE TEST TYPE: S
26 PAGE
27 *
28 *
29 * BREAKPOINTS OR SENSE SWITCHES CONTROL THE RUNNING OF THE PROGRAM:
30 *
31 * SWITCH RESET SET
32 * -----
33 *
34 * BP 1 RUN IN THE NORMAL MODE STOP AT THE END OF THE CURRENT
35 * SW 1 AS DETERMINED BY THE PASS AND RETURN TO KEY BOARD
36 * OTHER SWITCHES. CONTROL
37 *
38 *
39 * BP 2 CONTINUE TO RUN THE CYCLE TEST RUNS FROM READ TO
40 * SW 2 TEST MODE SELECTED BY PUNCH TO READ, ETC.
41 * SWITCH 3.
42 *
43 *
44 * BP 3 SELECTS PUNCH MODE SELECTS READ MODE.
45 * SW 3 (USED WHEN SWITCH 2 IS RESET OR WHEN STARTING THE TEST)
46 *
47 *
48 * BP 4 STOP AND TYPE DIAGNOSTIC DO NOT STOP OR TYPE ON ERRORS.

```

Table 4-4. 925/930 Computers, Extended Mode I/O Test Program (Sheet 2 of 20)

```

49 * SW 4 MESSAGES WHENEVER AN CONTINUE RUNNING REGARDLESS.
50 * ERROR OCCURS.
51 *
52 *
53 * NOTE: TO RUN CONTINUOUS FROM READ TO PUNCH THE OUTPUT OF THE PUNCH
54 * SHOULD BE FED INTO THE READER. A LOOP SHOULD BE USED TO RUN
55 * CONTINUOUS ON THE READER.
56 * PAGE
57 *
58 * THE FOLLOWING SECTION OF CODE DEFINES EITHER A 925/930 INSTRUCTION
59 * SET OR A 9300 INSTRUCTION SET. THIS IS DONE ON THE BASIS OF THE
60 * MACHINE TYPE STATEMENT AT THE BEGINNING OF THE DECK.
61 *
62 *
63 *
64 *
65 PROC
66 DF9300 NAME
67 $INST FORM 3,6,15
68 P PROC 1 DEFINE INDEX OP'S TO IMPLY X1.
69 $LDX NAME 017
70 $STX NAME 077
71 $BRX NAME 057
72 INST (P(*1)*2)+1,P(0),P(1)
73 END
74 P PROC 1
75 $XAB NAME 037733 DEFINE COMPATIBLE REGISTER OP'S
76 $ABC NAME 037731
77 $BAC NAME 037713
78 $CLR NAME 037711
79 INST 0,040,P(0)
80 END
81 P PROC 1
82 $RSH NAME 000 DEFINE COMPATIBLE SHIFTS
83 $RCY NAME 002
84 $LSH NAME 004
85 $LCY NAME 006
86 $NOD NAME 044
87 $SHFT FORM 3,6,6,9
88 SHFT P(2),060,P(0),P(1)**0777
89 END
90 P PROC 1 DEFINE MISC. COMPATIBLE OP'S.
91 $MIN NAME 071
92 $MIB NAME 030
93 $BIM NAME 032
94 INST (P(*1)*2)+P(2),P(0),P(1)
95 END
96 *

```


Table 4-4. 925/930 Computers, Extended Mode I/O Test Program (Sheet 3 of 20)

```

97 $X      EQU      1              DEFINE INDEX TAG *X* FOR X1
98 *
99          AORG     0177          EOD CONSTANT
100 $EODC  DATA    04000000      EOD CONSTANT
101 *
102 $$SWTFM FORM    3,6,3,12
103 Q       PROC     1
104 $BPT    NAME     040000
105 $BVT    NAME     014040
106 $ROV    NAME     04000
107 $$SOV   NAME     00040
108        SWTFM    0,022,Q(0)*/(-12)**7,Q(0)**07777++(Q(1)>0)*/(6-Q(1))
109        END
110        END
111 *
112        PROC
113 DF930   NAME
114 *
115 * IF NOT A 9300 THEN DO THE FOLLOWING OPERATIONS
116 *
117 P       PROC     1              DEFINE I/O INSTRUCTIONS
118 $MIB    NAME     012
119 $BIM    NAME     032
120 INST   FORM     3,6,1,14
121        INST     P(2),P(0),P(*1),P(1)
122        END
123 *
124 $X      EQU      2              DEFINE INDEX TAG *X* FOR 2 ON 930
125 *
126 *
127        AORG     0177
128 $EODC  DATA    04000000      EOD CONSTANT
129 *
130 *
131 N       PROC     1              SET OVERFLOW
132 $$SOV   NAME
133        BRR      $,4
134        END
135 *
136 M       PROC     1              BRANCH AND CLEAR INTERRUPT
137 $BRC    NAME
138        DS       M(*1),1,2
139        BRU     M(1)
140        BRU     *$+1
141        PZE     M(1)
142        END
143 P       PROC     1
144 $DSC    NAME     0

```

Table 4-4. 925/930 Computers, Extended Mode I/O Test Program (Sheet 4 of 20)

	145	\$ALC	NAME	050000	
	146	\$ASC	NAME	012000	
	147	\$TSP	NAME	014000	
	148	I	FORM	3,6,15	
	149	I		$P(1)**2,(P(1)**4)/3++2,P(0)++(P(1)**1)/6$	
	150		END		
	151	P	PR0C	1	
	152	\$CAT	NAME	014000	
	153	\$CET	NAME	011000	
	154	\$CIT	NAME	010400	
	155	\$CZT	NAME	012000	
	156	I	FORM	3,6,15	
	157	I		$P(1)**2,040,P(0)++(P(1)**4)/12++(P(1)**1)/6$	
	158		END		
	159		END		
	160		D0	TYPE=9300	
	161		DF9300		
	162		D0	TYPE<9300	
00177	00400000		DF930		
			164	PAGE	
			165	*	
			166	* THE START OF THE PROGRAM	
			167	*	
00200			168	AORG	0200
00200	0 76 0	00203	169	BEGIN LDA	BRUG0 INITIALIZE RECOVERY LOCATIONS
00201	0 35 0	00001	170	STA	1
00202	0 35 0	00032	171	STA	032
00203	0 01 0	00204	172	BRUG0 BRU	KYBD
00204	0 23 0	01241	173	KYBD EXU	RDIS
00205	0 23 0	01231	174	EXU	PDIS
00206	0 46 30003		175	G02 CLR	
00207	0 35 0	01444	176	STA	T1
00210	0 02 0	02001	177	RKB	0,1,1
00211	0 32 0	01445	178	G01 BIM	T1+1 READ CHARACTER
00212	0 76 0	01445	179	LDA	T1+1
00213	0 75 0	01636	180	LDB	=077
00214	0 70 0	01637	181	SKM	=° P° CHECK FOR CONTROL CHAR.
00215	0 01 0	00217	182	BRU	\$+2
00216	0 01 0	00233	183	BRU	P0
00217	0 70 0	01640	184	SKM	=° R°
00220	0 01 0	00222	185	BRU	\$+2
00221	0 01 0	00256	186	BRU	R0
00222	0 70 0	01641	187	SKM	=° S°
00223	0 01 0	00225	188	BRU	\$+2
00224	0 43 0	00317	189	BRM	S0
00225	0 72 0	01642	190	SKA	=070 CHECK FOR DIGIT
00226	0 01 0	00206	191	BRU	G02 IF NOT CONTROL OR DIGIT CLEAR
00227	0 75 0	01444	192	LDB	T1

Table 4-4. 925/930 Computers, Extended Mode I/O Test Program (Sheet 5 of 20)

00230	0	36	0	01446	193	STB	T1+2	SAVE CHANNEL NUMBER IN T1+2	
00231	0	35	0	01444	194	STA	T1	SAVE UNIT NUMBER IN T1	
00232	0	01	0	00211	195	BRU	G01		
					196	PAGE			
					197	*			
					198	*	PUNCH CHANNEL SET UP ROUTINE		
					199	*			
00233	0	43	0	00323	200	PO	BRM	MAKECH	GET CHANNEL NO. BUILT. SAVED IN T1+2
00234	0	71	0	01643	201		LDX	=-4**0177777	UPDATE EGM/EOD'S
00235	2	76	0	01234	202		LDA	PALC+1,X	
00236	0	14	0	01644	203		ETR	=050277677	
00237	0	16	0	01446	204		MRG	T1+2	INSERT CHANNEL DESIGNATION
00240	0	35	1	00235	205		STA	*\$-3	
00241	0	41	0	00235	206		BRX	\$-4	
00242	0	71	0	01645	207		LDX	=-3**0177777	UPDATE SKS'S
00243	2	76	0	01227	208		LDA	PCET+1,X	
00244	0	14	0	01646	209		ETR	=057737677	
00245	0	16	0	01447	210		MRG	T1+3	
00246	0	35	1	00243	211		STA	*\$-3	
00247	0	41	0	00243	212		BRX	\$-4	
00250	0	71	0	01643	213		LDX	=-4**0177777	UPDATE CHANNEL COMMANDS
00251	2	76	0	01212	214		LDA	PIOSD+2,X	
00252	0	14	0	01647	215		ETR	=050277777	
00253	0	16	0	01450	216		MRG	T1+4	
00254	0	35	1	00251	217		STA	*\$-3	
00255	0	41	0	00256	218		BRX	\$+1	
00256	0	41	0	00251	219		BRX	\$-5	
00257	0	76	0	01230	220		LDA	PTL	UPDATE UNIT NO.
00260	0	14	0	01650	221		ETR	=-2	
00261	0	16	0	01444	222		MRG	T1	
00262	0	35	0	01230	223		STA	PTL	
00263	0	16	0	01651	224		MRG	=02000	MAKE NO L. DER EGM
00264	0	35	0	01227	225		STA	PPT	
00265	0	01	0	00234	226		BRU	KYBD	
					227	*			
					228	*	READER CHANNEL SET UP ROUTINE		
					229	*			
					230	*			
00266	0	43	0	00323	231	RO	BRM	MAKECH	GET CHANNEL NO BUILT
00267	0	71	0	01643	232		LDX	=-4**0177777	UPDATE EGM/EOD'S
00270	2	76	0	01244	233		LDA	RALC+1,X	
00271	0	14	0	01644	234		ETR	=050277677	
00272	0	16	0	01446	235		MRG	T1+2	
00273	0	35	1	00270	236		STA	*\$-3	
00274	0	41	0	00270	237		BRX	\$-4	
00275	0	71	0	01643	238		LDX	=-4**0177777	BUILD SKS'S
00276	2	76	0	01240	239		LDA	RCIT+1,X	
00277	0	14	0	01646	240		ETR	=057737677	

Table 4-4. 925/930 Computers, Extended Mode I/O Test Program (Sheet 6 of 20)

00300	0	16	0	01447	241	MRG	T1+3	
00301	0	35	1	00276	242	STA	*S-3	
00302	0	41	0	00276	243	BRX	\$-4	
00303	0	71	0	01652	244	LDX	=-10**0177777	BUILD CHANNEL COMMAND EOM/EOD'S
00304	2	76	0	01224	245	LDA	R10RD+2,X	
00305	0	14	0	01653	246	ETR	=070277777	
00306	0	16	0	01450	247	MRG	T1+4	
00307	0	35	1	00304	248	STA	*S-3	
00310	0	41	0	00311	249	BRX	\$+1	
00311	0	41	0	00304	250	BRX	\$-5	
00312	0	76	0	01240	251	LDA	RPT	BUILD RPT WITH UNIT NO.
00313	0	14	0	01650	252	ETR	=-2	
00314	0	16	0	01444	253	MRG	T1	
00315	0	35	0	01240	254	STA	RPT	
00316	0	01	0	00204	255	BRU	KYBD	
					256	*		
					257	*		
					258	*	START TEST RUNNING	
					259	*		
00317	0	02	00000		260	SO DSC	0	DISCONNECT
00320	0	40	20100		261	BPT	3	START READ OR PUNCH?
00321	0	01	0 00640		262	BRU	IN	READ
00322	0	01	0 00356		263	BRU	OUT	PUNCH
					264		PAGE	
					265	*		
					266	*	BUILD CHANNEL NO. SUBROUTINE	
					267	*		
00323	0	00	0 00000		268	MAKECH	PZE	
00324	0	46	30003		269	CLR		
00325	0	76	0 01446		270	LDA	T1+2	GET CH. NO.
00326	0	72	0 01654		271	SKA	=4	EOD REQUIRED
00327	0	75	0 00177		272	LDB	E0DC	YES
00330	0	36	0 01450		273	STB	T1+4	NO, SAVE EOD BIT.
00331	0	14	0 01655		274	ETR	=3	
00332	0	35	0 01446		275	STA	T1+2	
00333	0	71	0 01446		276	LDX	T1+2	
00334	0	76	0 01450		277	LDA	T1+4	
00335	2	16	0 00352		278	MRG	MAKETB,X	BUILD EOM/EOD SELECTION
00336	0	35	0 01446		279	STA	T1+2	SAVE EOM
00337	0	46	10012		280	BAC		BUILD SKS SELECTION
00340	0	72	0 00177		281	SKA	E0DC	
00341	0	75	0 01656		282	LDB	=040000	
00342	0	46	00014		283	XAB		
00343	2	16	0 00352		284	MRG	MAKETB,X	
00344	0	35	0 01447		285	STA	T1+3	SAVE SKS
00345	0	76	0 01444		286	LDA	T1	BUILD UNIT NO. BIT
00346	0	14	0 01657		287	ETR	=1	
00347	0	17	0 01657		288	EOR	=1	

Table 4-4. 925/930 Computers, Extended Mode I/O Test Program (Sheet 7 of 20)

00350	0 35 0 01444	289	STA	T1	SAVE UNIT NO. BIT
00351	0 51 0 00323	290	BRR	MAKECH	EXIT
		291	*		
00352	00000000	292	MAKETB	DATA	0,0100,020000000,020000100
00353	00000100				
00354	20000000				
00355	20000100				
		293	PAGE		
		294	*		
		295	*	PAPER TAPE PUNCH OUTPUT SECTION.	
		296	*		
00356	0 71 0 01650	297	OUT	LDX	=-64**0177777 SET UP OUTPUT IMAGE WITH 64 WORDS
00357	0 46 30003	298		CLR	
00360	2 35 0 01344	299		STA	IMAGE+64,X
00361	0 55 0 01651	300		ADD	=01000000
00362	0 41 0 00350	301		BRX	\$-2
00363	0 23 0 01230	302	OUT4	EXU	PTL START PUNCH WITH LEADER
00364	0 23 0 01233	303		EXU	PALC ALER1 CHANNEL
00365	0 23 0 01206	304		LCH	PI0SP LOAD I0SP IMAGE,64
00366	0 13 0 01207				
00367	0 71 0 01643	305		LDX	=-4**0177777 SET UP FOR SUBROUTINE TO TYPE:
00370	0 75 0 01652	306		LDB	=ERMSG1 I0SP, OUTPUT
00371	0 76 0 00431	307		LDA	PCATC
00372	0 43 0 00610	308		BRM	WCZ WAIT FOR CHANNEL COUNT = ZERO
00373	0 23 0 01226	309		EXU	PCET ERROR ON OUTPUT
00374	0 40 20040	310		BPT	4 YES, CHECK BP4 FOR NO STOP
00375	0 01 0 00416	311		BRU	OUT4A NO, NO STOP ON ERROR
00376	0 23 0 01231	312	OUT2	EXU	PDIS ERROR, DISCONNECT CHANNEL
00377	0 02 0 02641	313		TYP	0,1,4 TYPE GENERAL ERROR MESSAGE
00400	0 12 0 01653	314		MIB	=052121225 I E
00401	0 12 0 01654	315		MIB	=051514651 RR9R
00402	0 12 0 01655	316		MIB	=012246451 DJR
00403	0 12 0 01656	317		MIB	=031452712 INS
00404	0 46 00014	318		XAB	TYPE SPECIFIC ERROR MESSAGE
00405	0 16 0 01635	319		MRG	MIBX
00406	0 35 0 00407	320		STA	\$+1
00407	0 12 0 00000	321		MIB	00 TO BE REPLACED AT RUN TIME
00410	0 61 0 00407	322		MIN	\$-1
00411	0 41 0 00407	323		BRX	\$-2
00412	0 02 14000	324	G0T0P	T0P	0
00413	0 40 14000	325		CAT	0
00414	0 01 0 00413	326		BRU	\$-1
00415	0 01 0 00204	327		BRU	KYBD
		328	*		
00416	0 43 0 00476	329	OUT4A	BRM	OUTPIN G3 PIN AND CHECK CHANNEL ADDRESS
00417	0 23 0 01224	330		EXU	PCAT IF CHANNEL INACTIVE READDRESS PUNCH
00420	0 01 0 00422	331		BRU	\$+2
00421	0 23 0 01227	332		EXU	PPT

Table 4-4. 925/930 Computers, Extended Mode I/O Test Program (Sheet 8 of 20)

00422	0 23 0 01233	333	EXU	PALC		
00423	0 23 0 01210	334	LCH	PI0SD		LOAD CH. I0SD IMAGE,64
00424	0 13 0 01211					
00425	0 71 0 01643	335	0UT1	LDX	=-4**0177777	SET UP ERROR MESSAGE:
00426	0 75 0 01657	336		LDB	=ERMSG2	I0SD, 0UTPUT
00427	0 76 0 00431	337		LDA	PCATC	
00430	0 43 0 00610	338		BRM	WCZ	WAIT FOR COUNT EQUAL ZERO
00431	0 23 0 01224	339	PCATC	EXU	PCAT	CHAN. ACTIVE?
00432	0 01 0 00431	340		BRU	\$-1	YES
00433	0 23 0 01226	341		EXU	PCET	NO, ERROR?
00434	0 40 20040	342		BPT	4	YES, IS STOP ALLOWED?
00435	0 01 0 00437	343		BRU	\$\$+2	NO ERROR OR NO STOP
00436	0 01 0 00376	344		BRU	0UT2	ERROR STOP ALLOWED
00437	0 43 0 00476	345		BRM	0UTPIN	GO PIN AND CHECK CHANNEL ADDRESS
00440	0 46 30003	346		CLR		
00441	0 35 0 01452	347		STA	PRF	R(PUNCH REPEAT)
00442	0 23 0 01230	348	0UT1A	EXU	PTL	
00443	0 23 0 01233	349		EXU	PALC	
00444	0 23 0 01210	350		LCH	PI0SD	LOAD I0SD IMAGE,64
00445	0 13 0 01211					
00446	0 71 0 01643	351		LDX	=-4**0177777	SET UP ERROR MESSAGE:
00447	0 75 0 01657	352		LDB	=ERMSG2	
00450	0 76 0 00431	353		LDA	PCATC	
00451	0 43 0 00610	354		BRM	WCZ	
00452	0 23 0 01224	355		EXU	PCAT	CHAN. ACTIVE?
00453	0 01 0 00452	356		BRU	\$-1	YES
00454	0 23 0 01226	357		EXU	PCET	NO, CHAN. ERROR
00455	0 40 20040	358		BPT	4	YES, IS ERROR STOP ALLOWED
00456	0 01 0 00450	359		BRU	\$\$+2	NO,NO
00457	0 01 0 00376	360		BRU	0UT2	YES
00460	0 43 0 00476	361		BRM	0UTPIN	GO PIN AND CHECK CHANNEL ADDRESS
00461	0 53 0 01452	362		SKN	PRF	PUNCH REPEAT
00462	0 01 0 00454	363		BRU	\$\$+2	RESET
00463	0 01 0 00457	364		BRU	0UT3	SET
00464	0 76 0 01670	365		LDA	\$\$-1	S(PUNCH REPEAT)
00465	0 35 0 01452	366		STA	PRF	
00466	0 01 0 00442	367		BRU	0UT1A	
		368	*			
		369	*			
00467	0 40 20400	370	0UT3	BPT	1	GO OR STOP
00470	0 01 0 00234	371		BRU	KYBD	STOP
00471	0 40 20200	372		BPT	2	GO, CYCLE?
00472	0 01 0 00640	373		BRU	IN	YES
00473	0 40 20100	374		BPT	3	ONE ONLY
00474	0 01 0 00640	375		BRU	IN	
00475	0 01 0 00356	376		BRU	0UT	
		377	*			
		378	*			SUBROUTINE TO CHECK CHANNEL ADDRESS

Table 4-4. 925/930 Computers, Extended Mode I/O Test Program (Sheet 9 of 20)

00476	0	00	0	00030	379 *			
00477	0	23	0	01232	380	OUTPIN	PZE	0
00500	0	33	0	01444	381		EXU	PASC
00501	0	76	0	01444	382		PIN	T1
00502	0	17	0	01671	383		LDA	T1
00503	0	72	0	01670	384		EOR	=IMAGE+64
00504	0	40	20040		385		SKA	=-1
00505	0	51	0	00476	386		BPT	4
00506	0	37	0	01445	387		BRR	OUTPIN
00507	0	46	00014		388		STX	T1+1
00510	0	16	0	01635	389		XAB	
00511	0	35	0	00530	390		MRG	MIBX
00512	0	76	0	01671	391		STA	OTPIN1
00513	0	43	0	00537	392		LDA	=IMAGE+64
00514	0	35	0	00577	393		BRM	MK9CT
00515	0	36	0	00630	394		STA	OTPNM1
00516	0	76	0	01444	395		STB	OTPNM1+1
00517	0	43	0	00537	396		LDA	T1
00520	0	35	0	00635	397		BRM	MK9CT
00521	0	36	0	00636	398		STA	OTPNM2
00522	0	23	0	01231	399		STB	OTPNM2+1
00523	0	71	0	01672	400		EXU	PDIS
00524	0	02	0	02641	401		LDX	=-15**017777
00525	2	12	0	00574	402		TYP	0,1,4
00526	0	41	0	00525	403		MIB	OTPNM+15,X
00527	0	71	0	01445	404		BRX	\$-1
00530	0	12	0	00030	405		LDX	T1+1
00531	0	61	0	00530	406	OTPIN1	MIB	00
00532	0	41	0	00530	407		MIN	\$-1
00533	0	71	0	01673	408		BRX	\$-2
00534	2	12	0	00610	409	OTPIN2	LDX	=-12**017777
00535	0	41	0	00534	410		MIB	OTPNM2+3,X
00536	0	01	0	00412	411		BRX	\$-1
					412		BRU	GETOP
					413 *			
					414 *			
					415 *	SUBROUTINE TO MAKE ONE WORD INTO 8 BCD OCTAL DIGITS		
					416 *			
00537	0	00	0	00030	417	MK9CT	PZE	0
00540	0	71	0	01674	418		LDX	=-8**017777
00541	0	46	20005		419		ABC	0 TO WORD TO 8
00542	0	6700	003		420		LSH	3
00543	2	35	0	01456	421		STA	T1+10,X
00544	0	76	0	01675	422		LDA	=0
00545	0	41	0	00542	423		BRX	\$-3
00546	0	46	30003		424		CLR	
00547	0	71	0	01674	425		LDX	=-8**017777
00550	0	6720	006		426		LCY	6
								REASSEMBLY BCD CHARACTERS INTO A + B.

Table 4-4. 925/930 Computers, Extended Mode I/O Test Program (Sheet 10 of 20)

00551	2 16 0 01456	427	MRC	T1+10,X	
00552	0 41 0 00550	428	BRX	\$-2	
00553	0 46 00014	429	XAB		
00554	0 51 0 00537	430	BRR	MK0CT	EXIT
		431	*		
		432	*		
		433	*	ERROR MESSAGES FOR CHANNEL ADDRESS TEST SUBROUTINE	
		434	*		
00555	52254524	435	0TPNM BCD	52, IEND ADDRESS FROM CHANNEL DID NOT AGREE WITH EXPECTE	
00556	12212424				
00557	51256262				
00560	12265146				
00561	44122330				
00562	21454525				
00563	43122431				
00564	24124546				
00565	63122127				
00566	51252512				
00567	66316330				
00570	12256747				
00571	25236325				
00572	24122126	436	BCD	8,D AFTER	
00573	63255112				
00574	25674725	437	BCD	12, EXPECTED	
00575	23632524				
00576	12121212				
00577	00000000	438	0TPNM1 DATA	0,0	
00600	00000000				
00601	73121252	439	BCD	16,, IRECEIVED	
00602	51252325				
00603	31652524				
00604	12121212				
00605	00000000	440	0TPNM2 DATA	0,0	
00606	00000000				
00607	33125252	441	BCD	4,, 11	
		442	PAGE		
		443	*		
		444	*	WAIT FOR COUNT EQUAL ZERO SUBROUTINE.	
		445	*		
00610	0 00 0 00000	446	WCZ PZE		ENTRY
00611	0 35 0 00617	447	STA	\$+6	SAVE R/PCAT
00612	0 55 0 01657	448	ADD	=1	MAKE A CZT
00613	0 35 0 00614	449	STA	\$+1	SAVE R/PCZT
00614	0 40 12000	450	CZT	00	C = 02?
00615	0 01 0 00617	451	BRU	\$+2	NO
00616	0 51 0 00610	452	BRR	WCZ	YES
00617	0 40 14000	453	CAT	00	CHANNEL ACTIVE?
00620	0 01 0 00614	454	BRU	\$-4	YES

Table 4-4. 925/930 Computers, Extended Mode I/O Test Program (Sheet 11 of 20)

00621	0	40	20040	455	BPT	4	NO, INHIBIT ERRORS
00622	0	51	0 00610	456	BRR	WCZ	YES
00623	0	37	0 01444	457	STX	T1	NO, PRINT ERROR MESSAGE
00624	0	46	00014	458	XAB		
00625	0	16	0 01635	459	MRG	MIBX	
00626	0	35	0 00634	460	STA	\$+6	
00627	0	71	0 01676	461	LDX	=-13**0177777	
00630	0	02	0 02641	462	TYP	0,1,4	
00631	2	12	0 01530	463	MIB	DISMSG+13,X	
00632	0	41	0 00631	464	BRX	\$-1	
00633	0	71	0 01444	465	LDX	T1	
00634	0	12	0 00030	466	MIE	00	
00635	0	61	0 00634	467	MIN	\$-1	
00636	0	41	0 00634	468	BRX	\$-2	
00637	0	01	0 00412	469	BRU	G0T0P	
				470		PAGE	
				471	*		
				472	*	INPUT SECTION	
				473	*		
00640	0	43	0 01177	474	IN	BRM	STARTP
00641	0	02	20001	475		R0V	START READER
00642	0	23	0 01243	476		EXU	RALC
00643	0	23	0 01212	477		LCH	R10SD
00644	0	13	0 01213				ALERT
00645	0	23	0 01235	478	IN0B	EXU	RCZT
00646	0	01	0 00657	479		BRU	INC
00647	0	23	0 01234	480	RCATC	EXU	RCAT
00650	0	40	20040	481		BPT	4
00651	0	01	0 00670	482		BRU	IN0A
00652	0	71	0 01676	483		LDX	=-13**0177777
00653	0	02	0 02641	484		TYP	0,1,4
00654	2	12	0 01525	485		MIB	ERMSG4+13,X
00655	0	41	0 00654	486		BRX	\$-1
00656	0	01	0 00412	487		BRU	G0T0P
				488	*		
00657	0	23	0 01234	489	IN0	EXU	RCAT
00660	0	01	0 00645	490		BRU	IN0B
00661	0	23	0 01235	491		EXU	RCZT
00662	0	01	0 00654	492		BRU	\$+2
00663	0	01	0 00670	493		BRU	IN0A
00664	0	71	0 01645	494		LDX	=-3**0177777
00665	0	75	0 01677	495		LDB	=ERMSG5
00666	0	76	0 00647	496		LDA	RCATC
00667	0	43	0 00610	497		BRM	WCZ
00670	0	76	0 01730	498	IN0A	LDA	=033120152
00671	0	75	0 01731	499		LDB	=BUFFER+64
00672	0	43	0 01055	500		BRM	CHECK
00673	0	43	0 01177	501	IN1	BRM	STARTP
							GO DO DISCONNECT ERROR TEST
							BL0CK NO. 1
							END ADDRESS EXPECTED
							GO CHECK DATA
							START READER IF DATA CHECKED O.K.

Table 4-4. 925/930 Computers, Extended Mode I/O Test Program (Sheet 12 of 20)

00674	0 02 20001	502	R0V		
00675	0 23 0 01243	503	EXU	RALC	
00676	0 23 0 01214	504	LCH	RI0SP	LOAD I0SP BUFFER,65
00677	0 13 0 01215				
00700	0 76 0 01670	505	LDA	=-1	
00701	0 35 0 01451	506	STA	SPF	S(I0SP INPUT)
00702	0 23 0 01237	507 IN4	EXU	RCIT	INTERRECORD CONDITION
00703	0 01 0 00735	508	BRU	\$+2	NO
00704	0 01 0 00732	509	BRU	IN2	YES
00705	0 23 0 01234	510	EXU	RCAT	CHAN. ACTIVE
00706	0 01 0 00716	511	BRU	IN1A	YES
00707	0 40 20040	512	BPT	4	NO, ERROR STOP PERMITTED
00710	0 01 0 00732	513	BRU	IN2	NO
00711	0 71 0 01672	514	LDX	=-15**0177777	YES
00712	0 02 0 02641	515	TYP	0,1,4	
00713	2 12 0 01547	516	MIB	ERMSG6+15,X	
00714	0 41 0 00713	517	BRX	\$-1	
00715	0 01 0 00412	518	BRU	GETOP	
		519 *			
00716	0 23 0 01235	520 IN1A	EXU	RCZT	C=0
00717	0 01 0 00732	521	BRU	IN4	NO
00720	0 23 0 01237	522	EXU	RCIT	YES, CHAN. INTER-RECORD
00721	0 01 0 00732	523	BRU	IN2	NO
00722	0 40 20040	524	BPT	4	YES, ERROR STOP PERMITTED
00723	0 01 0 00732	525	BRU	IN2	NO
00724	0 23 0 01241	526	EXU	RDIS	YES, STOP TAPE
00725	0 71 0 01732	527	LDX	=-17**0177777	
00726	0 02 0 02641	528	TYP	0,1,4	
00727	2 12 0 01570	529	MIB	ERMSG7+17,X	
00730	0 41 0 00727	530	BRX	\$-1	
00731	0 01 0 00412	531	BRU	GETOP	
		532 *			
00732	0 23 0 01236	533 IN2	EXU	RCET	IF CHANNEL ERROR
00733	4 51 0 00733	534	S0V		SET OVERFLOW
00734	0 23 0 01241	535	EXU	RDIS	STOP TAPE
00735	0 53 0 01451	536	SKN	SPF	I0SP INPUT FLAG
00736	0 01 0 00756	537	BRU	IN3	RESET
00737	0 76 0 01733	538	LDA	=033120252	SET, GET BLOCK NO. 2
00740	0 75 0 01444	539	LDB	BUFFER+64	END ADDRESS EXPECTED
00741	0 43 0 01055	540	BRM	CHECK	CHECK DATA INPUT
		541 *			IF CORRECT CONTINUE.
00742	0 43 0 01177	542	BRM	STARTP	START TAPE
00743	0 02 20001	543	R0V		
00744	0 23 0 01243	544	EXU	RALC	ALERT
00745	0 23 0 01216	545	LCH	RI0SP1	LOAD I0SP BUFFER,32
00746	0 13 0 01217				
00747	0 71 0 01645	546	LDX	=-3**0177777	SET UP ERROR MESSAGE
00750	0 75 0 01734	547	LDB	=ERMSG8	

Table 4-4. 925/930 Computers, Extended Mode I/O Test Program (Sheet 13 of 20)

00751	0	76	0	00647	548	LDA	RCATC	
00752	0	43	0	00610	549	BRM	WCZ	GO WAIT FOR C=0
00753	0	23	0	01236	550	EXU	RCET	IF CHANNEL ERROR
00754	4	51	0	00754	551	S0V		SET OVERFLOW
00755	0	23	0	01234	552	EXU	RCAT	IS READER STILL RUNNING
00756	0	01	0	00750	553	BRU	IN5	YES
00757	0	23	0	01240	554	EXU	RPT	NO, RESTART TAPE READER
00760	0	23	0	01243	555	EXU	RALC	
00761	0	23	0	01220	556	LCH	RI0RP	LOAD I0RD BUFFER+32,33
00762	0	13	0	01221				
00763	0	46	0	30003	557	CLR		
00764	0	35	0	01451	558	STA	SPF	R(I0SP INPUT)
00765	0	01	0	00732	559	BRU	IN4	
					560	*		
00766	0	76	0	01735	561	IN3	LDA	=033120352
00767	0	75	0	01731	562	LDB	=BUFFER+64	BLOCK NO. 3
00770	0	43	0	01055	563	BRM	CHECK	END ADDRESS EXPECTED
00771	0	43	0	01177	564	BRM	STARTP	GO CHECK DATA
00772	0	23	0	01243	565	EXU	RALC	START TAPE
00773	0	23	0	01222	566	LCH	RI0RD	LOAD I0RD BUFFER,56
00774	0	13	0	01223				
00775	0	23	0	01234	567	EXU	RCAT	CHAN. ACTIVE
00776	0	01	0	00775	568	BRU	\$-1	YES, WAIT FOR STOP
00777	0	71	0	01674	569	LDX	=-8**0177777	CHAN. INACTIVE
01000	0	76	0	01636	570	LDA	=077	
01001	2	72	0	01444	571	SKA	BUFFER+64,X	CHECK FOR 0'S IN LAST 8 CHARACTERS
01002	0	31	0	01024	572	BRU	IN3A	
01003	0	41	0	01031	573	BRX	\$-2	
01004	0	71	0	01674	574	IN3B	LDX	=-8**0177777
01005	0	76	0	01642	575	LDA	=070	O.K. INSERT CORRECT 8 CHARACTERS
01006	2	35	0	01444	576	STA	BUFFER+64,X	
01007	0	55	0	01657	577	ADD	=1	
01010	0	41	0	01036	578	BRX	\$-2	
01011	0	76	0	01736	579	LDA	=033120452	BLOCK NO. 4
01012	0	75	0	01737	580	LDB	=BUFFER+56	END ADDRESS EXPECTED
01013	0	02	0	20001	581	R0V		
01014	0	43	0	01055	582	BRM	CHECK	GO CHECK DATA
01015	0	40	0	20400	583	BPT	1	TEST STOP?
01016	0	01	0	00234	584	BRU	KYBD	YES
01017	0	40	0	20200	585	BPT	2	TEST CYCLE?
01020	0	01	0	00356	586	BRU	0UT	YES
01021	0	40	0	20100	587	BPT	3	TEST ONE ONLY.
01022	0	01	0	00640	588	BRU	IN	READ
01023	0	01	0	00356	589	BRU	0UT	PUNCH
					590	*		
01024	0	40	0	20040	591	IN3A	BPT	4
01025	0	01	0	01034	592	BRU	IN3B	NO
01026	0	71	0	01674	593	LDX	=-8**0177777	FORMAT LAST EIGHT CHARACTERS FOR TYPE

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01027	0	46	30003	594	CLR		
01030	2	76	0 01444	595	IN3D	LDA	BUFFER+64,X
01031	0	6700	003	596		LSH	3
01032	2	16	0 01444	597		MRG	BUFFER+64,X
01033	0	14	0 01710	598		ETR	=0707
01034	0	16	0 01711	599		MRG	=052120000
01035	0	6700	006	600		LSH	6
01036	2	35	0 01055	601		STA	MSGIMG+8,X
01037	0	41	0 01030	602		BRX	IN3D
01040	0	71	0 01712	603		LDX	=-19**0177777 OUTPUT ERROR MESSAGE
01041	0	02	0 02641	604		TYP	0,1,4
01042	2	12	0 01616	605		MIB	ERMSG9+19,X
01043	0	41	0 01042	606		BRX	\$-1
01044	0	02	14000	607		TOP	0
01045	0	40	14000	608		CAT	0
01046	0	01	0 01045	609		BRU	\$-1
01047	0	02	0 02041	610		TYP	0,1,1 OUTPUT 8 CHARACTERS
01050	0	71	0 01674	611		LDX	=-8**0177777
01051	2	12	0 01055	612		MIB	MSGIMG+8,X
01052	0	41	0 01051	613		BRX	\$-1
01053	0	12	0 01533	614		MIB	ERMSG1+3 CR
01054	0	01	0 00412	615		BRU	GETOP
				616	*		
				617	*		
01055				618		MSGIMG RES	8
				619		PAGE	
				620	*		
				621	*	CHECK INPUT DATA SUBROUTINE.	
				622	*		
01065	0	00	0 00030	623	CHECK	PZE	
01066	0	35	0 01626	624		STA	ERMSG0+6 SAVE BLOCK NO.
01067	0	23	0 01236	625		EXU	RCET CHECK FOR ERROR
01070	0	01	0 01141	626		BRU	PARERR GO TO PARITY ERROR ROUTINE
01071	0	40	20001	627		0VT	CHECK FOR PREVIOUSLY NOTED ERROR
01072	0	01	0 01141	628		BRU	PARERR GO TO PARITY ERROR ROUTINE
01073	0	23	0 01242	629		EXU	RASC STORE CHANNEL ADDRESS
01074	0	33	0 01444	630		PIN	T1
01075	0	36	0 01445	631		STB	T1+1 SAVE EXPECTED
01076	0	76	0 01444	632		LDA	T1
01077	0	17	0 01445	633		EOR	T1+1 COMPARE ACTUAL WITH EXPECTED
01100	0	72	0 01670	634		SKA	=-1 AGREE
01101	0	01	0 01154	635		BRU	PINERR NO
01102	0	71	0 01650	636	CHECK2	LDX	=-64**0177777 YES
01103	0	76	0 01675	637		LDA	=0C
01104	0	75	0 01636	638		LDB	=077
01105	2	70	0 01444	639		SKM	BUFFER+64,X CHECK INPUT BUFFER
01106	0	01	0 01112	640		BRU	CHECK1 ERROR
01107	0	55	0 01657	641		ADD	=1

Table 4-4. 925/930 Computers, Extended Mode I/O Test Program (Sheet 15 of 20)

01110	0 41 0 01105	642	BRX	\$-3	
01111	0 51 0 01055	643	BRR	CHECK	EXIT IF ALL CORRECT
		644	*		
01112	0 40 20040	645	CHECK1 BPT	4	ERROR STOP PERMITTED?
01113	0 51 0 01055	646	BRR	CHECK	NO, EXIT
01114	0 35 0 01444	647	STA	T1	YES
01115	0 6700 003	648	LSH	3	FORMAT EXPECTED
01116	0 16 0 01444	649	MRG	T1	
01117	0 14 0 01710	650	ETR	=0707	
01120	0 6700 006	651	LSH	6	
01121	0 16 0 01713	652	MRG	=012000052	
01122	0 35 0 01631	653	STA	ERMSG0+9	STORE EXPECTED
01123	2 76 0 01444	654	LDA	BUFFER+64,X	FORMAT RECEIVED
01124	0 5600 003	655	RSH	3	
01125	2 76 0 01444	656	LDA	BUFFER+64,X	
01126	0 6700 003	657	LSH	3	
01127	0 14 0 01710	658	ETR	=0707	
01130	0 6700 006	659	LSH	6	
01131	0 16 0 01713	660	MRG	=012000052	
01132	0 35 0 01634	661	STA	ERMSG0+12	STORE RECEIVED
01133	0 71 0 01676	662	LDX	=-13**0177777	
01134	0 23 0 01241	663	EXU	RDIS	DISCONNECT READER CHANNEL
01135	0 02 0 02641	664	TYP	0,1,4	
01136	2 12 0 01635	665	MIB	ERMSG0+13,X	
01137	0 41 0 01136	666	BRX	\$-1	
01140	0 01 0 00412	667	BRU	G0T0P	
		668	*		
		669	* PARITY ERROR SUBROUTINE		
		670	*		
01141	0 40 20040	671	PARERR BPT	4	ERROR STOP PERMITTED?
01142	0 51 0 01055	672	BRR	CHECK	NO
01143	0 23 0 01241	673	EXU	RDIS	YES, DISCONNECT READER CHANNEL
01144	0 02 0 02641	674	TYP	0,1,4	
01145	0 71 0 01714	675	LDX	=-9**0177777	
01146	2 12 0 01627	676	MIB	ERMSGP+9,X	OUTPUT PARITY ERROR MESSAGE
01147	0 41 0 01146	677	BRX	\$-1	
01150	0 02 14000	678	T0P	0	
01151	0 40 14000	679	CAT	0	
01152	0 01 0 01151	680	BRU	\$-1	
01153	0 01 0 01102	681	BRU	CHECK2	RETURN TO CHECK NUMBERS
		682	*		
		683	* PIN ADDRESS ERROR SUBROUTINE		
		684	*		
01154	0 40 20040	685	PINERR BPT	4	ERROR STOP PERMITTED
01155	0 01 0 01102	686	BRU	CHECK2	NO
01156	0 76 0 01444	687	LDA	T1	YES
01157	0 43 0 00537	688	BRM	MK0CT	EXPAND ACTUAL TO BCD
01160	0 35 0 00635	689	STA	0TPNM2	SAVE

Table 4-4. 925/930 Computers, Extended Mode I/O Test Program (Sheet 16 of 20)

01161	0 36 0 00636	690	STB	0TPNM2+1	
01162	0 76 0 01445	691	LDA	T1+1	EXPAND EXPECTED TO BCD
01163	0 43 0 00537	692	BRM	MK0CT	
01164	0 35 0 00577	693	STA	0TPNM1	SAVE
01165	0 36 0 00630	694	STB	0TPNM1+1	
01166	0 23 0 01241	695	EXU	RDIS	DISCONNECT READER
01167	0 02 0 02641	696	TYP	0,1,4	
01170	0 71 0 01672	697	LDX	=-15**0177777	
01171	2 12 0 00574	698	MIB	0TPNM+15,X	OUTPUT GENERAL MESSAGE
01172	0 41 0 01171	699	BRX	\$-1	
01173	0 12 0 01624	700	MIB	ERMSG0+4	OUTPUT BLOCK NO.
01174	0 12 0 01625	701	MIB	ERMSG0+5	
01175	0 12 0 01626	702	MIB	ERMSG0+6	
01176	0 01 0 00533	703	BRU	0TPIN2	GO OUTPUT RECEIVED AND EXPECTED
		704	*		
		705	*	START TAPE READER SUBROUTINE	
		706	*		
01177	0 00 0 00030	707	STARTP	PZE	
01200	0 71 0 01650	708	LDX	=-64**0177777	
01201	0 46 30003	709	CLR		
01202	2 35 0 01444	710	STA	BUFFER+64,X	CLEAR BUFFER
01203	0 41 0 01232	711	BRX	\$-1	
01204	0 23 0 01240	712	EXU	RPT	START TAPE
01205	0 51 0 01177	713	BRK	STARTP	EXIT
		714		PAGE	
		715	*		
		716	*	I/O CHANNEL COMMANDS	
		717	*		
01206	002 146 0 30	718	PI0SP	I0SP	IMAGE,64
01207	0100 01244				
01210	002 142 0 30	719	PI0SD	I0SD	IMAGE,64
01211	0100 01244				
		720	*		
01212	002 142 0 30	721	RI0SD	I0SD	BUFFER,64
01213	0100 01344				
01214	002 146 0 30	722	RI0SP	I0SP	BUFFER,65
01215	0101 01344				
01216	002 146 0 30	723	RI0SP1	I0SP	BUFFER,32
01217	0040 01344				
01220	002 144 0 30	724	RI0RP	I0RP	BUFFER+32,33
01221	0041 01404				
01222	002 140 0 30	725	RI0RD	I0RD	BUFFER,56
01223	0070 01344				
		726	*		
		727	*		
		728	*	I/O CHANNEL INSTRUCTIONS.	
		729	*		
01224	0 40 14000	730	PCAT	CAT	0

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01225	0 40 12000	731	PCZT	CZT	0	
01226	0 40 11000	732	PCET	CET	0	
		733	*			
01227	0 02 0 02044	734	PPT	PPT	0,1,1	
01230	0 02 0 00044	735	PTL	PTL	0,1,1	
01231	0 02 00000	736	PDIS	DSC	0	
01232	0 02 12000	737	PASC	ASC	0	
01233	0 02 50000	738	PALC	ALC	0	
		739	*			
		740	*			
01234	0 40 14000	741	RCAT	CAT	0	
01235	0 40 12000	742	RCZT	CZT	0	
01236	0 40 11000	743	RCET	CET	0	
01237	0 40 10400	744	RCIT	CIT	0	
		745	*			
01240	0 02 0 02034	746	RPT	RPT	0,1,1	
01241	0 02 00000	747	RDIS	DSC	0	
01242	0 02 12000	748	RASC	ASC	0	
01243	0 02 50000	749	RALC	ALC	0	
		750	*			
		751	*			
		752	*	OUTPUT IMAGE AREA, INPUT BUFFER AREA		
		753	*			
01244		754	IMAGE	RES	64	
01344		755	BUFFER	RES	64	
		756	*			
		757	*			
		758	*	TEMPORARY STORAGE AND FLAGS		
		759	*			
01444		760	T1	RES	13	
01461		761	SPF	RES	1	IOSP INPUT FLAG
01462		762	PRF	RES	1	PUNCH REPEAT FLAG
		763		PAGE		
		764	*			
		765	*	ERROR AND STATUS MESSAGES.		
		766	*			
01463	52233021	767	DISMSG	BCD	52	1CHANNEL ERROROUSLY DISCONNECTED BEFORE C=0, DURING
01464	45452543					
01465	12255151					
01466	46514664					
01467	62437012					
01470	24316223					
01471	46454525					
01472	23632524					
01473	12222526					
01474	46512512					
01475	23130073					
01476	12246451					

Table 4-4. 925/930 Computers, Extended Mode I/O Test Program (Sheet 18 of 20)

01477	31452712		
01500	31466247	768 ERMSG1 BCD	16, IOSP, OUTPUT. 11
01501	73124664		
01502	63476463		
01503	33125252		
01504	31466224	769 ERMSG2 BCD	16, IOSD, OUTPUT. 11
01505	73124664		
01506	63476463		
01507	33125252		
01510	52233021	770 ERMSG4 BCD	52, CHANNEL DID NOT DISCONNECT WHEN C=0 ON IOSD INPUT!
01511	45452543		
01512	12243124		
01513	12454663		
01514	12243162		
01515	23464545		
01516	25236312		
01517	66302545		
01520	12231300		
01521	12464512		
01522	31466224		
01523	12314547		
01524	64635212		
01525	31466224	771 ERMSG5 BCD	12, IOSD, INPUT!
01526	73123145		
01527	47646352		
01530	52233021	772 ERMSG6 BCD	52, CHANNEL DISCONNECTED DURING IOSP INPUT, CIT NEVER
01531	45452543		
01532	12243162		
01533	23464545		
01534	25236325		
01535	24122464		
01536	51314527		
01537	12314662		
01540	47123145		
01541	47646373		
01542	12233163		
01543	12452565		
01544	25511212		
01545	63516425	773	BCD 8, TRUE. 1
01546	33125212		
01547	52246451	774 ERMSG7 BCD	48, DURING IOSP INPUT C=0 INDICATING EOR PAST BUT C
01550	31452712		
01551	31466247		
01552	12314547		
01553	64631223		
01554	13001231		
01555	45243123		
01556	21633145		

Table 4-4. 925/930 Computers, Extended Mode I/O Test Program (Sheet 19 of 20)

01557	27122546			
01560	51124721			
01561	62631222			
01562	64631223			
01563	31631266	775	BCD	20, IT WAS NEVER TRUE. I
01564	21621245			
01565	25652551			
01566	12635164			
01567	25331252			
01570	31466247	776	ERMSG8 BCD	12, I0SP, INPUT!
01571	73123145			
01572	47646352			
01573	52314651	777	ERMSG9 BCD	52, I10RD 0N INPUT DID N0I IGN0RE THE LAST 8 CHARACTERS.
01574	24124645			
01575	12314547			
01576	64631224			
01577	31241245			
01600	46631231			
01601	27454651			
01602	25126330			
01603	25124321			
01604	62631210			
01605	12233021			
01606	51212363			
01607	25516273			
01610	12226463	778	BCD	24, BUT READ THE FOLLOWING!
01611	12512521			
01612	24126330			
01613	25122646			
01614	43434666			
01615	31452752			
01616	52524721	779	ERMSGP BCD	8, I1PARITY
01617	51316370			
01620	52314547	780	ERMSG0 BCD	40, I1INPUT ERR0R IN BLOCK N0. N!EXPECTED DD!
01621	64631225			
01622	51514651			
01623	12314512			
01624	22434623			
01625	42124546			
01626	33124552			
01627	25674725			
01630	23632524			
01631	12242452			
01632	51252325	781	BCD	12, RECEIVED DD!
01633	31652524			
01634	12242452			
01635	0 12 0 0000	782 *		
		783 MIBX	MIB	0
				OUTPUT INSTRUCTION

Table 4-4. 925/930 Computers, Extended Mode I/O Test Program (Sheet 20 of 20)

		784 *		
		785 *		
		786	END	BEGIN
01636	00000200			
01636	00000077			
01637	00606047			
01640	00606051			
01641	00606062			
01642	00000070			
01643	00177774			
01644	50277677			
01645	00177775			
01646	57737677			
01647	50277777			
01650	77777776			
01651	00002000			
01652	00177766			
01653	70277777			
01654	00000004			
01655	00000003			
01656	00040000			
01657	00000001			
01660	00177700			
01661	01000000			
01662	00001500			
01663	52121225			
01664	51514651			
01665	12246451			
01666	31452712			
01667	00001504			
01670	77777777			
01671	00001344			
01672	00177761			
01673	00177764			
01674	00177770			
01675	00000000			
01676	00177763			
01677	00001525			
01700	33120152			
01701	00001444			
01702	00177757			
01703	33120252			
01704	00001570			
01705	33120352			
01706	33120452			
01707	00001434			
01710	00000707			
01711	52120000			
01712	00177755			
01713	12000052			
01714	00177767			

4.64 Operation

4.65 To select the reader and punch units to be used enter:

"CU" P to select punch,

"CU" R to select reader.

The letter C is the channel number and may be any digit 0-7 and the letter U is the unit, either one or two. The reader and punch need not be on the same channel.

4.66 The test is started by entering the letter "S". Control of the test operation is then a function of the Breakpoint Switches. Table 4-5 summarizes the switch functions.

4.67 When running cyclic tests from punch to read, the tape from the punch should be inserted into the reader.

4.68 Punching

4.69 The program punches four blocks of 64 characters each in one pass. The characters form a counting sequence from 00 to 77g. The first block is started with leader and output with an IOSP. All punching is done in the one character per word mode. When the word count reaches zero, an IOSD is loaded to punch a second block of 64 characters. No leader is punched between the first and second blocks. This results in one physical block 128 characters long. Starting with leader, two additional blocks of 64 characters are then punched with an IOSD.

4.70 At the conclusion of each output operation, the channel address register is stored and compared with the expected value. If they do not agree, the program types the expected and actual values.

4.71 The program tests the channel during the output operation to see if the channel should erroneously disconnect before the word count reaches zero.

4.72 Reading

4.73 Each of the four blocks is read with a different set of commands and counts so as to test as many operations as possible. After reading a block, a general subroutine checks for input parity errors, for channel end address agreement with the expected address, and the data read character by character. Error messages with block numbers are typed in the event of any one of these tests failing. In the event of a test failing, reference should then be made to the test program flow diagram and troubleshooting information contained in Section 5 of this manual. The handling of each block is as follows:

a. Block 1. The first block is one-half of a 128 character physical block. Reading one character per word, an IOSD with a count of 64 is used to read this block. The program checks to see if the count reaches zero and the channel becomes inactive at the same time.

b. Block 2. This is the second half of the first physical block and is read with an IOSP with a count of 65. The read should terminate because of the end of record. The program checks to see that the word count does not reach zero and the channel remains active after the CIT (inter-record test) instruction skips. The tape is

Table 4-5. 9300 Computer Breakpoint Switch Functions

Breakpoint Switch	Reset	Set
1	Run in the normal mode as determined by the other switches	Stop and return to keyboard control at the end of the current pass (punch or read)
2	Continue to run test selected by switch 3	Cycle test runs from punch to read to punch and so on
3*	Selects punch mode	Selects to punch and so on
4	Stop and type diagnostic messages whenever an error occurs	Do not stop and type on errors but continue to run

*Used when switch 2 is reset or when starting test

finally stopped with a disconnect before the data is checked.

c. Block 3. Block 3 is a 64 character physical block. It is read with two channel commands. The first is an IOSB with a count of 32. If the count goes to zero before the channel disconnects, an IORP with a count of 33 is loaded. This should cause the inter-record indicator to be turned on at the end of the record. The count should not reach zero and the channel should remain active. The tape is again stopped with a disconnect before the data is checked.

d. Block 4. This is the third physical block of 64 characters and is read with an IORD with a count of 56. The program waits for the channel to be inactive then checks to see if the channel ignored the last eight characters. If the tape was erroneously stopped after the 56th character, it will show up as a failure on the first block of the next read pass.

4.74 Test Program

4.75 Table 4-6 gives the test program for the 9300 computer I/O extended mode.

Table 4-6. 9300 Computer, Extended Mode I/O Test Program (Sheet 1 of 20)

00022124

```

1 TYPE EQU 9300 MACHINE DEFINITION: 9300
2 *
3 *
4 * EXTENDED MODE I/O TEST PROGRAM. (925/930/9300 ALL CHANNELS)
5 *
6 * A.W. ENGLAND, SDS
7 *
8 *
9 * THIS PROGRAM USES THE PAPER TAPE READER AND PUNCH TO CHECK THE
10 * OPERATION OF THE I/O CHANNELS IN THE EXTENDED INTERLACE MODE OF
11 * OPERATION. THE READER AND PUNCH MAY BE CONNECTED TO ANY INTERLACED
12 * IMCC OR DACC. THEY NEED NOT BE ATTACHED TO THE SAME CHANNEL. THE
13 * PROGRAM ADDRESSES THE KEYBOARD FOR INFORMATION ABOUT UNIT AND CHANNEL
14 * SELECTION.
15 *
16 * TO SELECT PUNCH TYPE: 'CU'P
17 *
18 * TO SELECT READER TYPE: 'CU'R
19 *
20 * WHERE 'C' REPRESENTS CHANNEL NO. AND MAY HAVE THE VALUES 0-7;
21 * AND 'U' REPRESENTS UNIT NO. AND MAY HAVE THE VALUE 1-2. THE LETTER
22 * P OR R CAUSE THE SELECTION TO BE MADE WITH THE TWO PREVIOUS DIGITS.
23 * SPACES SHOULD NOT BE TYPED BETWEEN DIGITS OR CONTROL CHARACTER.
24 *
25 * TO START THE TEST TYPE: S
26 PAGE
27 *
28 *
29 * BREAKPOINTS OR SENSE SWITCHES CONTROL THE RUNNING OF THE PROGRAM:
30 *
31 * SWITCH RESET SET
32 * -----
33 *
34 * BP 1 RUN IN THE NORMAL MODE STOP AT THE END OF THE CURRENT
35 * SW 1 AS DETERMINED BY THE PASS AND RETURN TO KEY BOARD
36 * OTHER SWITCHES. CONTROL
37 *
38 *
39 * BP 2 CONTINUE TO RUN THE CYCLE TEST RUNS FROM READ TO
40 * SW 2 TEST MODE SELECTED BY PUNCH TO READ, ETC.
41 * SWITCH 3.
42 *
43 *
44 * BP 3 SELECTS PUNCH MODE SELECTS READ MODE.
45 * SW 3 (USED WHEN SWITCH 2 IS RESET OR WHEN STARTING THE TEST)
46 *
47 *
48 * BP 4 STOP AND TYPE DIAGNOSTIC DO NOT STOP OR TYPE ON ERRORS.

```

Table 4-6. 9300 Computer, Extended Mode I/O Test Program (Sheet 2 of 20)

```

49 *      SW 4      MESSAGES WHENEVER AN      CONTINUE RUNNING REGARDLESS.
50 *      ERROR OCCURS.
51 *
52 *
53 * NOTE: TO RUN CONTINUOUS FROM READ TO PUNCH THE OUTPUT OF THE PUNCH
54 *      SHOULD BE FED INTO THE READER. A LOOP SHOULD BE USED TO RUN
55 *      CONTINUOUS ON THE READER.
56 *      PAGE
57 *
58 * THE FOLLOWING SECTION OF CODE DEFINES EITHER A 925/930 INSTRUCTION
59 * SET OR A 9300 INSTRUCTION SET. THIS IS DONE ON THE BASIS OF THE
60 * MACHINE TYPE STATEMENT AT THE BEGINNING OF THE DECK.
61 *
62 *
63 *
64 *
65 *      PROC
66 DF9300 NAME
67 $INST FORM      3,6,15
68 P      PROC      1                      DEFINE INDEX OP'S TO IMPLY X1.
69 $LDX  NAME      017
70 $STX  NAME      077
71 $BRX  NAME      057
72      INST      (P(*1)*2)++1,P(0),P(1)
73      END
74 P      PROC      1
75 $XAB  NAME      037733                   DEFINE COMPATIBLE REGISTER OP'S
76 $ABC  NAME      037731
77 $BAC  NAME      037713
78 $CLR  NAME      037711
79      INST      0,040,P(0)
80      END
81 P      PROC      1
82 $RSH  NAME      000                      DEFINE COMPATIBLE SHIFTS
83 $RCY  NAME      002
84 $LSH  NAME      004
85 $LCY  NAME      006
86 $NOD  NAME      044
87 $SHFT FORM      3,6,6,9
88      SHFT      P(2),060,P(0),P(1)**0777
89      END
90 P      PROC      1                      DEFINE MISC. COMPATIBLE OP'S.
91 $MIN  NAME      071
92 $MIB  NAME      030
93 $RIM  NAME      032
94      INST      (P(*1)*2)++P(2),P(0),P(1)
95      END
96 *

```

Table 4-6. 9300 Computer, Extended Mode I/O Test Program (Sheet 3 of 20)

```

97 $X      EQU      1              DEFINE INDEX TAG 'X' FOR X1
98 *
99          AORG     0177          EOD CONSTANT
100 $E9DC   DATA    04000000     EOD CONSTANT
101 *
102 $$SWFRM FORM     3,6,3,12
103 Q       PROC     1
104 $BPT    NAME     040000
105 $QVT    NAME     014040
106 $R8V    NAME     04000
107 $$QV    NAME     00040
108        SWFRM    0,022,Q(0)*7/(-12)**7,Q(0)**07777777+(Q(1)>0)*7/(5-Q(1))
109        END
110        END
111 *
112        PROC
113 DF930   NAME
114 *
115 * IF NOT A 9300 THEN DO THE FOLLOWING OPERATIONS
116 *
117 P       PROC     1              DEFINE I/O INSTRUCTIONS
118 $MIB    NAME     012
119 $BIM    NAME     032
120 INST   FORM     3,6,1,14
121        INST     P(2),P(0),P(*1),P(1)
122        END
123 *
124 $X      EQU      2              DEFINE INDEX TAG 'X' FOR 2 ON 930
125 *
126 *
127          AORG     0177
128 $E9DC   DATA    04000000     EOD CONSTANT
129 *
130 *
131 N       PROC     1              SET OVERFLOW
132 $$QV    NAME
133        BRR      $,4
134        END
135 *
136 M       PROC     1              BRANCH AND CLEAR INTERRUPT
137 $BRC    NAME
138        DB       M(*1),1,2
139        BRU      M(1)
140        BRU      *$+1
141        PZE      M(1)
142        END
143 P       PROC     1
144 $DSC    NAME     0

```

Table 4-6. 9300 Computer, Extended Mode I/O Test Program (Sheet 4 of 20)

```

145 $ALC NAME 050000
146 $ASC NAME 012000
147 $TOP NAME 014000
148 I FORM 3,6,15
149 I P(1)**2,(P(1)**4)/3++2,P(0)+(P(1)**1)*/6
150 END
151 P PRGC 1
152 $CAT NAME 014000
153 $CET NAME 011000
154 $CIT NAME 010400
155 $CZT NAME 012000
156 I FORM 3,6,15
157 I P(1)**2,040,P(0)+(P(1)**4)/12+(P(1)**1)*/6
158 END
159 END
160 D0 TYPE=9300
00177 04000000 161 DF9300
162 D0 TYPE<9300
163 DF930
164 PAGE
165 *
166 * THE START OF THE PROGRAM
167 *
00200 168 AORG 0200
00200 0 0 16 00203 169 BEGIN LDA BRUG0 INITIALIZE RECOVERY LOCATIONS
00201 0 0 76 00031 170 STA 1
00202 0 0 76 00032 171 STA 032
00203 0 0 01 00204 172 BRUG0 BRU KYBD
00204 0 0 21 01241 173 KYBD EXU RDIS
00205 0 0 21 01231 174 EXU PDIS
00206 0 40 37711 175 G02 CLR
00207 0 0 76 01444 176 STA T1
00210 0 02 0 02031 177 RKB 0,1,1
00211 0 32 01445 178 G01 BIM T1+1 READ CHARACTER
00212 0 0 16 01445 179 LDA T1+1
00213 0 0 14 01636 180 LDB =077
00214 0 0 55 01637 181 SKM =' P' CHECK FOR CONTROL CHAR.
00215 0 0 01 00217 182 BRU $+2
00216 0 0 01 00233 183 BRU P0
00217 0 0 55 01640 184 SKM =' R'
00220 0 0 01 00222 185 BRU $+2
00221 0 0 01 00256 186 BRU R0
00222 0 0 55 01641 187 SKM =' S'
00223 0 0 01 00225 188 BRU $+2
00224 0 0 03 00317 189 BRM S0
00225 0 0 54 01642 190 SKA =070 CHECK FOR DIGIT
00226 0 0 01 00206 191 BRU G02 IF NOT CONTROL OR DIGIT CLEAR
00227 0 0 14 01444 192 LDB T1

```


Table 4-6. 9300 Computer, Extended Mode I/O Test Program (Sheet 5 of 20)

00230	0 0 74	01446	193	STB	T1+2	SAVE CHANNEL NUMBER IN T1+2
00231	0 0 76	01444	194	STA	T1	SAVE UNIT NUMBER IN T1
00232	0 0 01	00211	195	BRU	G01	
			196	PAGE		
			197	*		
			198	*	PUNCH CHANNEL SET UP ROUTINE	
			199	*		
00233	0 0 03	00323	200	PO BRM	MAKECH	GET CHANNEL NO. BUILT. SAVED IN T1+2
00234	1 17	01643	201	LDX	=-4**0177777	UPDATE E0M/E0D'S
00235	0 1 16	01234	202	LDA	PALC+1,X	
00236	0 0 11	01644	203	ETR	=050277677	
00237	0 0 13	01446	204	MRG	T1+2	INSERT CHANNEL DESIGNATION
00240	1 0 76	00235	205	STA	*\$-3	
00241	1 57	00235	206	BRX	\$-4	
00242	1 17	01645	207	LDX	=-3**0177777	UPDATE SKS'S
00243	0 1 16	01227	208	LDA	PCFT+1,X	
00244	0 0 11	01646	209	ETR	=057737677	
00245	0 0 13	01447	210	MRG	T1+3	
00246	1 0 76	00243	211	STA	*\$-3	
00247	1 57	00243	212	BRX	\$-4	
00250	1 17	01643	213	LDX	=-4**0177777	UPDATE CHANNEL COMMANDS
00251	0 1 16	01212	214	LDA	PJ0SD+2,X	
00252	0 0 11	01647	215	ETR	=050277777	
00253	0 0 13	01450	216	MRG	T1+4	
00254	1 0 76	00251	217	STA	*\$-3	
00255	1 57	00256	218	BRX	\$+1	
00256	1 57	00251	219	BRX	\$-5	
00257	0 0 16	01230	220	LDA	PTL	UPDATE UNIT NO.
00260	0 0 11	01650	221	ETR	=-2	
00261	0 0 13	01444	222	MRG	T1	
00262	0 0 76	01230	223	STA	PTL	
00263	0 0 13	01651	224	MRG	=02000	MAKE NO LEADER E0M
00264	0 0 76	01227	225	STA	PPT	
00265	0 0 01	00234	226	BRU	KYRD	
			227	*		
			228	*	READER CHANNEL SET UP ROUTINE	
			229	*		
			230	*		
00266	0 0 03	00323	231	RC BRM	MAKECH	GET CHANNEL NO BUILT
00267	1 17	01643	232	LDX	=-4**0177777	UPDATE E0M/E0D'S
00270	0 1 16	01244	233	LDA	RALC+1,X	
00271	0 0 11	01644	234	ETR	=050277677	
00272	0 0 13	01446	235	MRG	T1+2	
00273	1 0 76	00270	236	STA	*\$-3	
00274	1 57	00270	237	BRX	\$-4	
00275	1 17	01643	238	LDX	=-4**0177777	BUILD SKS'S
00276	0 1 16	01240	239	LDA	RCIT+1,X	
00277	0 0 11	01646	240	ETR	=057737677	

Table 4-6. 9300 Computer, Extended Mode I/O Test Program (Sheet 6 of 20)

00300	0 0 13 01447	241	MRG	T1+3	
00301	1 0 76 00276	242	STA	*\$-3	
00302	1 57 00276	243	BRX	\$-4	
00303	1 17 01652	244	LDX	=-10**0177777	BUILD CHANNEL COMMAND EOM/EOD'S
00304	0 1 16 01224	245	LDA	RIORD+2,X	
00305	0 0 11 01653	246	ETR	=070277777	
00306	0 0 13 01450	247	MRG	T1+4	
00307	1 0 76 00304	248	STA	*\$-3	
00310	1 57 00311	249	BRX	\$+1	
00311	1 57 00304	250	BRX	\$-5	
00312	0 0 16 01240	251	LDA	RPT	BUILD RPT WITH UNIT NO.
00313	0 0 11 01650	252	ETR	=-2	
00314	0 0 13 01444	253	MRG	T1	
00315	0 0 76 01240	254	STA	RPT	
00316	0 0 01 00204	255	BRU	KYBD	
		256	*		
		257	*		
		258	*	START TEST RUNNING	
		259	*		
00317	0 02 00000	260	SO DSC	0	DISCONNECT
00320	0 22 4 0010	261	BPT	3	START READ OR PUNCH?
00321	0 0 01 00640	262	BRU	IN	READ
00322	0 0 01 00356	263	BRU	OUT	PUNCH
		264		PAGE	
		265	*		
		266	*	BUILD CHANNEL NO. SUBROUTINE	
		267	*		
00323	0 0 00 00000	268	MAKECH	PZE	
00324	0 40 37711	269		CLR	
00325	0 0 16 01446	270	LDA	T1+2	GET CH. NO.
00326	0 0 54 01654	271	SKA	=4	EOD REQUIRED
00327	0 0 14 00177	272	LDB	EODC	YES
00330	0 0 74 01450	273	STB	T1+4	NO, SAVE EOD BIT.
00331	0 0 11 01655	274	ETR	=3	
00332	0 0 76 01446	275	STA	T1+2	
00333	1 17 01446	276	LDX	T1+2	
00334	0 0 16 01450	277	LDA	T1+4	
00335	0 1 13 00352	278	MRG	MAKETB,X	BUILD EOM/EOD SELECTION
00336	0 0 76 01446	279	STA	T1+2	SAVE EOM
00337	8 40 37713	280	BAC		BUILD SKS SELECTION
00340	0 0 54 00177	281	SKA	EODC	
00341	0 0 14 01656	282	LDB	=040000	
00342	0 40 37733	283	XAB		
00343	0 1 13 00352	284	MRG	MAKETB,X	
00344	0 0 76 01447	285	STA	T1+3	SAVE SKS
00345	0 0 16 01444	286	LDA	T1	BUILD UNIT NO. BIT
00346	0 0 11 01657	287	ETR	=1	
00347	0 0 12 01657	288	EOR	=1	

Table 4-6. 9300 Computer, Extended Mode I/O Test Program (Sheet 7 of 20)

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00350	0 0 76 01444	289	STA	T1	SAVE UNIT NO. BIT
00351	0 0 41 00323	290	BRR	MAKECH	EXIT
		291	*		
00352	00000000	292	MAKETB	DATA	0,0100,020000000,020000100
00353	00000100				
00354	20000000				
00355	20000100				
		293		PAGE	
		294	*		
		295	*	PAPER TAPE PUNCH OUTPUT SECTION.	
		296	*		
00356	1 17 01660	297	OUT	LDX	=-64**0177777 SET UP OUTPUT IMAGE WITH 64 WORDS
00357	0 40 37711	298		CLR	
00360	8 1 76 01344	299		STA	IMAGE+64,X
00361	0 0 05 01651	300		ADD	=01000000
00362	1 57 00360	301		BRX	\$-2
00363	0 0 21 01230	302	OUT4	EXU	PTL START PUNCH WITH LEADER
00364	0 0 21 01233	303		EXU	PALC ALERT CHANNEL
00365	0 0 21 01236	304		LCH	PIOSP LOAD IOSP IMAGE,64
00366	0 0 31 01237				
00367	1 17 01643	305		LDX	=-4**0177777 SET UP FOR SUBROUTINE TO TYPE:
00370	8 0 14 01652	306		LDB	=ERMSG1 IOSP, OUTPUT
00371	0 0 16 00431	307		LDA	PCATC
00372	0 0 03 00610	308		BRM	WCZ WAIT FOR CHANNEL COUNT = ZERO
00373	0 0 21 01226	309		EXU	PCET ERROR ON OUTPUT
00374	0 22 4 0004	310		BPT	4 YES, CHECK BP4 FOR NO STOP
00375	0 0 01 00416	311		BRU	OUT4A NO, NO STOP ON ERROR
00376	0 0 21 01231	312	OUT2	EXU	PDIS ERROR, DISCONNECT CHANNEL
00377	0 02 0 02641	313		TYP	0,1,4 TYPE GENERAL ERROR MESSAGE
00400	0 30 01663	314		MIB	=052121225 I E
00401	0 30 01664	315		MIB	=051514651 RROR
00402	0 30 01665	316		MIB	=012246451 DUR
00403	0 30 01666	317		MIB	=031452712 ING
00404	0 40 37733	318		XAB	TYPE SPECIFIC ERROR MESSAGE
00405	0 0 13 01635	319		MRG	MIBX
00406	0 0 76 00437	320		STA	\$+1
00407	0 30 00000	321		MIB	00 TO BE REPLACED AT RUN TIME
00410	0 71 00407	322		MIN	\$-1
00411	1 57 00407	323		BRX	\$-2
00412	0 02 14000	324	GETOP	TOP	0
00413	0 20 14000	325		CAT	0
00414	0 0 01 00413	326		BRU	\$-1
00415	0 0 31 00234	327		BRU	KYED
		328	*		
00416	0 0 03 00476	329	OUT4A	BRM	OUTPIN GO PIN AND CHECK CHANNEL ADDRESS
00417	0 0 21 01224	330		EXU	PCAT IF CHANNEL INACTIVE READDRESS PUNCH
00420	0 0 01 00422	331		BRU	\$+2
00421	0 0 21 01227	332		EXU	PPT

SDS 900685C

Table 4-6. 9300 Computer, Extended Mode I/O Test Program (Sheet 8 of 20)

00422	0 0 21 01233	333	EXU	PALC		
00423	0 0 21 01210	334	LCH	PI0SD	LOAD CH. I0SD IMAGE,64	
00424	0 0 31 01211					
00425	1 17 01643	335	OUT1	LDX	=-4**0177777	SET UP ERROR MESSAGE:
00426	0 0 14 01657	336	LDB	=ERMSG2	I0SD, OUTPUT	
00427	0 0 16 00431	337	LDA	PCATC		
00430	0 0 03 00610	338	BRM	WCZ	WAIT FOR COUNT EQUAL ZERO	
00431	0 0 21 01224	339	PCATC	EXU	PCAT	CHAN. ACTIVE?
00432	0 0 01 00431	340	BRU	\$-1	YES	
00433	0 0 21 01226	341	EXU	PCET	NO, ERROR?	
00434	0 22 4 0004	342	BPT	4	YES, IS STOP ALLOWED?	
00435	0 0 01 00437	343	BRU	\$+2	NO ERROR OR NO STOP	
00436	0 0 01 00376	344	BRU	OUT2	ERROR STOP ALLOWED	
00437	0 0 03 00476	345	BRM	OUTPIN	GO PIN AND CHECK CHANNEL ADDRESS	
00440	0 40 37711	346	CLR			
00441	0 0 76 01452	347	STA	PRF	R(PUNCH REPEAT)	
00442	0 0 21 01230	348	OUT1A	EXU	PTL	
00443	0 0 21 01233	349	EXU	PALC		
00444	0 0 21 01210	350	LCH	PI0SD	LOAD I0SD IMAGE,64	
00445	0 0 31 01211					
00446	1 17 01643	351	LDX	=-4**0177777	SET UP ERROR MESSAGE:	
00447	0 0 14 01657	352	LDB	=ERMSG2		
00450	0 0 16 00431	353	LDA	PCATC		
00451	0 0 03 00610	354	BRM	WCZ		
00452	0 0 21 01224	355	EXU	PCAT	CHAN. ACTIVE?	
00453	0 0 01 00452	356	BRU	\$-1	YES	
00454	0 0 21 01226	357	EXU	PCET	NO, CHAN. ERROR	
00455	0 22 4 0004	358	BPT	4	YES, IS ERROR STOP ALLOWED	
00456	0 0 01 00450	359	BRU	\$+2	NO, NO	
00457	0 0 01 00376	360	BRU	OUT2	YES	
00460	0 0 03 00476	361	BRM	OUTPIN	GO PIN AND CHECK CHANNEL ADDRESS	
00461	0 0 53 01452	362	SKN	PRF	PUNCH REPEAT	
00462	0 0 01 00454	363	BRU	\$+2	RESET	
00463	0 0 01 00457	364	BRU	OUT3	SET	
00464	0 0 16 01670	365	LDA	=-1	S(PUNCH REPEAT)	
00465	0 0 76 01452	366	STA	PRF		
00466	0 0 01 00442	367	BRU	OUT1A		
		368	*			
		369	*			
00467	0 22 4 0040	370	OUT3	BPT	1	GO OR STOP
00470	0 0 01 00204	371	BRU	KYBD		STOP
00471	0 22 4 0020	372	BPT	2		GO, CYCLE?
00472	0 0 01 00640	373	BRU	IN		YES
00473	0 22 4 0010	374	BPT	3		ONE ONLY
00474	0 0 01 00640	375	BRU	IN		
00475	0 0 01 00356	376	BRU	OUT		
		377	*			
		378	*	SUBROUTINE TO CHECK CHANNEL ADDRESS		

Table 4-6. 9300 Computer, Extended Mode I/O Test Program (Sheet 9 of 20)

00476	0 0 00 00000	379 *			
00477	0 0 21 01232	380	OUTPIN PZE	0	
00500	0 0 33 01444	381	EXU	PASC	ALERT TO STORE PUNCH CHANNEL
00501	0 0 16 01444	382	PIN	T1	STORE ADDRESS
00502	0 0 12 01671	383	LDA	T1	
00503	0 0 54 01670	384	EOR	=IMAGE+64	COMPARE WITH EXPECTED
00504	0 22 4 0004	385	SKA	=-1	
00505	0 0 41 00476	386	BPT	4	NOT EQUAL
00506	1 77 01445	387	BRR	OUTPIN	EQUAL OR ERROR STOP NOT PERMITTED
00507	0 40 37733	388	STX	T1+1	SAVE LENGTH OF ERROR MESSAGE
00510	0 0 13 01635	389	XAB		SET UP ERROR OUTPUTTER
00511	0 0 76 00530	390	MRG	MIBX	
00512	0 0 16 01671	391	STA	OTPIN1	
00513	0 0 03 00537	392	LDA	=IMAGE+64	GENERATE EXPECTED PIN WORD IN BCD
00514	0 0 76 00577	393	BRM	MK0CT	
00515	0 0 74 00630	394	STA	OTPNM1	SAVE EXPECTED
00516	0 0 16 01444	395	STB	OTPNM1+1	
00517	0 0 03 00537	396	LDA	T1	GENERATE ACTUAL PIN WORD IN BCD
00520	0 0 76 00635	397	BRM	MK0CT	
00521	0 0 74 00636	398	STA	OTPNM2	
00522	0 0 21 01231	399	STB	OTPNM2+1	
00523	1 17 01672	400	EXU	PDIS	DISCONNECT PUNCH CHANNEL
00524	0 02 0 02641	401	LDX	=-15**0177777	
00525	1 30 00574	402	TYP	0,1,4	
00526	1 57 00525	403	MIB	OTPNM+15,X	OUTPUT MESSAGE
00527	1 17 01445	404	BRX	\$-1	
00530	0 30 00000	405	LDX	T1+1	OUTPUT SPECIAL MESSAGE
00531	0 71 00530	406	OTPIN1 MIB	00	
00532	1 57 00530	407	MIN	\$-1	
00533	1 17 01673	408	BRX	\$-2	
00534	1 30 00610	409	OTPIN2 LDX	=-12**0177777	OUTPUT RECEIVED AND EXPECTED MESSAGE
00535	1 57 00534	410	MIB	OTPNM2+3,X	
00536	0 0 01 00412	411	BRX	\$-1	
		412	BRU	G0T0P	
		413 *			
		414 *			
		415 *	SUBROUTINE TO MAKE ONE WORD INTO 8 BCD OCTAL DIGITS		
		416 *			
00537	0 0 00 00000	417	MK0CT PZE	0	
00540	1 17 01674	418	LDX	=-8**0177777	
00541	0 40 37731	419	ABC		0 TO WORD TO B
00542	0 60 04 003	420	LSH	3	SHIFT OUT OCTAL DIGIT
00543	0 1 76 01456	421	STA	T1+10,X	SAVE BCD CHARACTER
00544	0 0 16 01675	422	LDA	=0	CLEAR A
00545	1 57 00542	423	BRX	\$-3	
00546	0 40 37711	424	CLR		
00547	1 17 01674	425	LDX	=-8**0177777	REASSEMBLY BCD CHARACTERS INTO A + B.
00550	0 60 06 005	426	LCY	6	

Table 4-6. 9300 Computer, Extended Mode I/O Test Program (Sheet 10 of 20)

00551	0 1 13 01456	427	MRG	T1+10,X	
00552	1 57 00550	428	BRX	\$-2	
00553	0 40 37733	429	XAB		
00554	0 0 41 00537	430	BRR	MK0CT	EXIT
		431	*		
		432	*		
		433	*	ERROR MESSAGES FOR CHANNEL ADDRESS TEST SUBROUTINE	
		434	*		
00555	52254524	435	0TPNM BCD	52,1END ADDRESS FROM CHANNEL DID NOT AGREE WITH EXPECTE	
00556	12212424				
00557	51256262				
00560	12265146				
00561	44122330				
00562	21454525				
00563	43122431				
00564	24124546				
00565	63122127				
00566	51252512				
00567	66316330				
00570	12256747				
00571	25236325				
00572	24122126	436	BCD	8,D AFTER	
00573	63255112				
00574	25674725	437	BCD	12,EXPECTED	
00575	23632524				
00576	12121212				
00577	00000000	438	0TPNM1 DATA	0,0	
00600	00000000				
00601	73121252	439	BCD	16,, IRECEIVED	
00602	51252325				
00603	31652524				
00604	12121212				
00605	00000000	440	0TPNM2 DATA	0,0	
00606	00000000				
00607	33125252	441	BCD	4,, 11	
		442	PAGE		
		443	*		
		444	*	WAIT FOR COUNT EQUAL ZERO SUBROUTINE.	
		445	*		
00610	0 0 00 00000	446	WCZ	PZE	ENTRY
00611	0 0 76 00617	447	STA	\$+6	SAVE R/PCAT
00612	0 0 05 01657	448	ADD	=1	MAKE A CZT
00613	0 0 76 00614	449	STA	\$+1	SAVE R/PCZT
00614	0 20 12000	450	CZT	00	C = 02?
00615	0 0 01 00617	451	BRU	\$+2	NO
00616	0 0 41 00610	452	BRR	WCZ	YES
00617	0 20 14000	453	CAT	00	CHANNEL ACTIVE?
00620	0 0 01 00614	454	BRU	\$-4	YES

Table 4-6. 9300 Computer, Extended Mode I/O Test Program (Sheet 11 of 20)

00621	0 22 4 0004	455	BPT	4	NO, INHIBIT ERRORS
00622	0 0 41 00610	456	BRR	WCZ	YES
00623	1 77 01444	457	STX	T1	NO, PRINT ERROR MESSAGE
00624	0 40 37733	458	XAB		
00625	0 0 13 01635	459	MRG	MIBX	
00626	0 0 76 00634	460	STA	\$+6	
00627	1 17 01676	461	LDX	=-13**0177777	
00630	0 02 0 02641	462	TYP	0.1.4	
00631	1 30 01500	463	MIB	DISMSG+13.X	
00632	1 57 00631	464	BRX	\$-1	
00633	1 17 01444	465	LDX	T1	
00634	0 30 00000	466	MIB	00	
00635	0 71 00634	467	MIN	\$-1	
00636	1 57 00634	468	BRX	\$-2	
00637	0 0 01 00412	469	BRU	GOTOP	
		470		PAGE	
		471 *			
		472 *	INPUT SECTION		
		473 *			
00640	0 0 03 01177	474 IN	BRM	STARTP	START READER
00641	0 22 0 4003	475	ROV		
00642	0 0 21 01243	476	EXU	RALC	ALERT
00643	0 0 21 01212	477	LCH	RIOSD	LOAD IOSD BUFFER,64
00644	0 0 31 01213				
00645	0 0 21 01235	478 INOB	EXU	RCZT	C=0
00646	0 0 01 00657	479	BRU	INO	NO
00647	0 0 21 01234	480 RCATC	EXU	RCAT	YES, CHAN. ACTIVE
00650	0 22 4 0004	481	BPT	4	YES, ERROR STOP PERMITTED
00651	0 0 01 00670	482	BRU	INOA	NO, NO CONT.
00652	1 17 01676	483	LDX	=-13**0177777	YES
00653	0 02 0 02641	484	TYP	0.1.4	
00654	1 30 01525	485	MIB	ERMSG4+13.X	
00655	1 57 00654	486	BRX	\$-1	
00656	0 0 01 00412	487	BRU	GOTOP	
		488 *			
00657	0 0 21 01234	489 INO	EXU	RCAT	CHAN. ACTIVE STILL?
00660	0 0 01 00645	490	BRU	INOB	YES
00661	0 0 21 01235	491	EXU	RCZT	NO, C=0
00662	0 0 01 00654	492	BRU	\$+2	NO
00663	0 0 01 00670	493	BRU	INOA	YES
00664	1 17 01645	494	LDX	=-3**0177777	SET UP ERROR MESSAGE
00665	0 0 14 01677	495	LDB	=ERMSG5	
00666	0 0 16 00647	496	LDA	RCATC	
00667	0 0 03 00610	497	BRM	WCZ	GO DO DISCONNECT ERROR TEST
00670	0 0 16 01700	498 INOA	LDA	=033120152	BLOCK NO. 1
00671	0 0 14 01701	499	LDB	=BUFFER+64	END ADDRESS EXPECTED
00672	0 0 03 01055	500	BRM	CHECK	GO CHECK DATA
00673	0 0 03 01177	501 IN1	BRM	STARTP	START READER IF DATA CHECKED O.K.

Table 4-6. 9300 Computer, Extended Mode I/O Test Program (Sheet 12 of 20)

00674	0 22 0 4003	502	R0V		
00675	0 0 21 01243	503	EXU	RALC	
00676	0 0 21 01214	504	LCH	RI0SP	LOAD I0SP BUFFER,65
00677	0 0 31 01215				
00700	0 0 16 01670	505	LDA	=-1	
00701	0 0 76 01451	506	STA	SPF	S(I0SP INPUT)
00702	0 0 21 01237	507 IN4	EXU	RCIT	INTERRECORD CONDITION
00703	0 0 01 00735	508	BRU	\$+2	NO
00704	0 0 01 00732	509	BRU	IN2	YES
00705	0 0 21 01234	510	EXU	RCAT	CHAN. ACTIVE
00706	0 0 01 00716	511	BRU	IN1A	YES
00707	0 22 4 0004	512	BPT	4	NO, ERROR STOP PERMITTED
00710	0 0 01 00732	513	BRU	IN2	NO
00711	1 17 01672	514	LDX	=-15**0177777	YES
00712	0 02 0 02641	515	TYP	0,1,4	
00713	1 30 01547	516	MIB	ERMSG6+15,X	
00714	1 57 00713	517	BRX	\$-1	
00715	0 0 01 00412	518	BRU	G0T0P	
		519 *			
00716	0 0 21 01235	520 IN1A	EXU	RC7T	C=0
00717	0 0 01 00732	521	BRU	IN4	NO
00720	0 0 21 01237	522	EXU	RCIT	YES, CHAN. INTER-RECORD
00721	0 0 01 00732	523	BRU	IN2	NO
00722	0 22 4 0004	524	BPT	4	YES, ERROR STOP PERMITTED
00723	0 0 01 00732	525	BRU	IN2	NO
00724	0 0 21 01241	526	EXU	RDIS	YES, STOP TAPE
00725	1 17 01702	527	LDX	=-17**0177777	
00726	0 02 0 02641	528	TYP	0,1,4	
00727	1 30 01570	529	MIB	ERMSG7+17,X	
00730	1 57 00727	530	BRX	\$-1	
00731	0 0 01 00412	531	BRU	G0T0P	
		532 *			
00732	0 0 21 01236	533 IN2	EXU	RCET	IF CHANNEL ERROR
00733	0 22 0 0043	534	S0V		SET OVERFLOW
00734	0 0 21 01241	535	EXU	RDIS	STOP TAPE
00735	0 0 53 01451	536	SKN	SPF	I0SP INPUT FLAG
00736	0 0 01 00756	537	BRU	IN3	RESEI
00737	0 0 16 01733	538	LDA	=033120252	SET, GET BLOCK NO. 2
00740	0 0 14 01444	539	LDB	BUFFER+64	END ADDRESS EXPECTED
00741	0 0 03 01055	540	BRM	CHECK	CHECK DATA INPUT
		541 *			IF CORRECT CONTINUE.
00742	0 0 03 01177	542	BRM	STARTP	START TAPE
00743	0 22 0 4003	543	R0V		
00744	0 0 21 01243	544	EXU	RALC	ALERT
00745	0 0 21 01216	545	LCH	RI0SP1	LOAD I0SP BUFFER,32
00746	0 0 31 01217				
00747	1 17 01645	546	LDX	=-3**0177777	SET UP ERROR MESSAGE
00750	0 0 14 01734	547	LDB	=ERMSG8	

Table 4-6. 9300 Computer, Extended Mode I/O Test Program (Sheet 13 of 20)

00751	0 0 16 00647	548	LDA	RCATC	
00752	0 0 03 00610	549	BRM	WCZ	GO WAIT FOR C=0
00753	0 0 21 01236	550	EXU	RCET	IF CHANNEL ERROR
00754	0 22 0 0043	551	S0V		SET OVERFLOW
00755	0 0 21 01234	552	EXU	RCAT	IS READER STILL RUNNING
00756	0 0 01 00750	553	BRU	INS	YES
00757	0 0 21 01240	554	EXU	RPT	NO, RESTART TAPE READER
00760	0 0 21 01243	555	INS EXU	RALC	
00761	0 0 21 01220	556	LCH	RI0RP	LOAD I0RD BUFFER+32,33
00762	0 0 31 01221				
00763	0 40 37711	557	CLR		
00764	0 0 76 01451	558	STA	SPF	R(I0SP INPUT)
00765	0 0 01 00732	559	BRU	IN4	
		560 *			
00766	0 0 16 01735	561	INS LDA	=033120352	BLOCK NO. 3
00767	0 0 14 01731	562	LDB	=BUFFER+64	END ADDRESS EXPECTED
00770	0 0 03 01055	563	BRM	CHECK	GO CHECK DATA
00771	0 0 03 01177	564	BRM	STARTP	START TAPE
00772	0 0 21 01243	565	EXU	RALC	
00773	0 0 21 01222	566	LCH	RI0RD	LOAD I0RD BUFFER,56
00774	0 0 31 01223				
00775	0 0 21 01234	567	EXU	RCAT	CHAN. ACTIVE
00776	0 0 01 00775	568	BRU	\$-1	YES, WAIT FOR STOP
00777	1 17 01674	569	LDX	=-8**0177777	CHAN. INACTIVE
01000	0 0 16 01636	570	LDA	=077	
01001	0 1 54 01444	571	SKA	BUFFER+64,X	CHECK FOR 0'S IN LAST 8 CHARACTERS
01002	0 0 01 01024	572	BRU	IN3A	
01003	1 57 01001	573	BRX	\$-2	
01004	1 17 01674	574	INS LDX	=-8**0177777	O.K. INSERT CORRECT 8 CHARACTERS
01005	0 0 16 01642	575	LDA	=070	
01006	0 1 76 01444	576	STA	BUFFER+64,X	
01007	0 0 05 01657	577	ADD	=1	
01010	1 57 01006	578	BRX	\$-2	
01011	0 0 16 01736	579	LDA	=033120452	BLOCK NO. 4
01012	0 0 14 01737	580	LDB	=BUFFER+56	END ADDRESS EXPECTED
01013	0 22 0 4003	581	R0V		
01014	0 0 03 01055	582	BRM	CHECK	GO CHECK DATA
01015	0 22 4 0043	583	BPT	1	TEST STOP?
01016	0 0 01 00234	584	BRU	KYBD	YES
01017	0 22 4 0023	585	BPT	2	TEST CYCLE?
01020	0 0 01 00356	586	BRU	0UT	YES
01021	0 22 4 0013	587	BPT	3	TEST ONE ONLY.
01022	0 0 01 00640	588	BRU	IN	READ
01023	0 0 01 00356	589	BRU	0UT	PUNCH
		590 *			
01024	0 22 4 0004	591	INS BPT	4	ERROR STOP PERMITTED
01025	0 0 01 01034	592	BRU	IN3B	NO
01026	1 17 01674	593	LDX	=-8**0177777	FORMAT LAST EIGHT CHARACTERS FOR TYPE

Table 4-6. 9300 Computer, Extended Mode I/O Test Program (Sheet 14 of 20)

01027	0 40 37711	594	CLR		
01030	0 1 16 01444	595	IN3D LDA	BUFFER+64,X	
01031	0 60 04 003	596	LSH	3	
01032	0 1 13 01444	597	MRG	BUFFER+64,X	
01033	0 0 11 01710	598	ETR	=0707	
01034	0 0 13 01711	599	MRG	=052120000	
01035	0 60 04 005	600	LSH	6	
01036	0 1 76 01055	601	STA	MSGIMG+8,X	
01037	1 57 01030	602	BRX	IN3D	
01040	1 17 01712	603	LDX	=-19**0177777	OUTPUT ERROR MESSAGE
01041	0 02 0 02641	604	TYP	0,1,4	
01042	1 30 01616	605	MIB	ERMSG9+19,X	
01043	1 57 01042	606	BRX	\$-1	
01044	0 02 14000	607	TSP	0	
01045	0 20 14000	608	CAT	0	
01046	0 0 01 01045	609	BRU	\$-1	
01047	0 02 0 02041	610	TYP	0,1,1	OUTPUT 8 CHARACTERS
01050	1 17 01674	611	LDX	=-8**0177777	
01051	1 30 01065	612	MIB	MSGIMG+8,X	
01052	1 57 01051	613	BRX	\$-1	
01053	0 30 01503	614	MIB	ERMSG1+3	CR
01054	0 0 01 09412	615	BRU	G0T0P	
		616	*		
		617	*		
01055		618	MSGIMG RES	8	
		619	PAGE		
		620	*		
		621	* CHECK INPUT DATA SUBROUTINE.		
		622	*		
01065	0 0 00 00000	623	CHECK PZE		
01066	0 0 76 01626	624	STA	ERMSG0+6	SAVE BLOCK NO.
01067	0 0 21 01236	625	EXU	RCET	CHECK FOR ERROR
01070	0 0 01 01141	626	BRU	PARERR	GO TO PARITY ERROR ROUTINE
01071	0 22 1 4040	627	0VT		CHECK FOR PREVIOUSLY NOTED ERROR
01072	0 0 01 01141	628	BRU	PARERR	GO TO PARITY ERROR ROUTINE
01073	0 0 21 01242	629	EXU	RASC	STORE CHANNEL ADDRESS
01074	0 0 33 01444	630	PIN	T1	
01075	0 0 74 01445	631	STB	T1+1	SAVE EXPECTED
01076	0 0 16 01444	632	LDA	T1	
01077	0 0 12 01445	633	EOR	T1+1	COMPARE ACTUAL WITH EXPECTED
01100	0 0 54 01670	634	SKA	=-1	AGREE
01101	0 0 01 01154	635	BRU	PINERR	NO
01102	1 17 01660	636	CHECK2 LDX	=-64**0177777	YES
01103	0 0 16 01675	637	LDA	=00	
01104	0 0 14 01636	638	LDB	=077	
01105	0 1 55 01444	639	SKM	BUFFER+64,X	CHECK INPUT BUFFER
01106	0 0 01 01112	640	BRU	CHECK1	ERROR
01107	0 0 05 01657	641	ADD	=1	
01110	1 57 01105	642	BRX	\$-3	

Table 4-6. 9300 Computer, Extended Mode I/O Test Program (Sheet 15 of 20)

01111	0 0 41 01055	643	BRR	CHECK	EXIT IF ALL CORRECT
		644	*		
01112	0 22 4 0004	645	CHECK1 BPT	4	ERROR STOP PERMITTED?
01113	0 0 41 01055	646	BRR	CHECK	NO, EXIT
01114	0 0 76 01444	647	STA	T1	YES
01115	0 60 04 003	648	LSH	3	FORMAT EXPECTED
01116	0 0 13 01444	649	MRG	T1	
01117	0 0 11 01710	650	ETR	=0707	
01120	0 60 04 005	651	LSH	6	
01121	0 0 13 01713	652	MRG	=012000052	
01122	0 0 76 01631	653	STA	ERMSG0+9	STORE EXPECTED
01123	0 1 16 01444	654	LDA	BUFFER+64,X	FORMAT RECEIVED
01124	0 50 00 003	655	RSH	3	
01125	0 1 16 01444	656	LDA	BUFFER+64,X	
01126	0 60 04 003	657	LSH	3	
01127	0 0 11 01710	658	ETR	=0707	
01130	0 60 04 005	659	LSH	6	
01131	0 0 13 01713	660	MRG	=012000052	
01132	0 0 76 01634	661	STA	ERMSG0+12	STORE RECEIVED
01133	1 17 01676	662	LDX	=-13**0177777	
01134	0 0 21 01241	663	EXU	RDIS	DISCONNECT READER CHANNEL
01135	0 02 0 02641	664	TYP	0,1,4	
01136	1 30 01635	665	MIB	ERMSG0+13,X	
01137	1 57 01136	666	BRX	\$-1	
01140	0 0 01 00412	667	BRU	G0T0P	
		668	*		
		669	*	PARITY ERROR SUBROUTINE	
		670	*		
01141	0 22 4 0004	671	PARERR BPT	4	ERROR STOP PERMITTED?
01142	0 0 41 01055	672	BRR	CHECK	NO
01143	0 0 21 01241	673	EXU	RDIS	YES, DISCONNECT READER CHANNEL
01144	0 02 0 02641	674	TYP	0,1,4	
01145	1 17 01714	675	LDX	=-9**0177777	
01146	1 30 01627	676	MIB	ERMSG0+9,X	OUTPUT PARITY ERROR MESSAGE
01147	1 57 01146	677	BRX	\$-1	
01150	0 02 14000	678	T0P	0	
01151	0 20 14000	679	CAT	0	
01152	0 0 01 01151	680	BRU	\$-1	
01153	0 0 01 01102	681	BRU	CHECK2	RETURN TO CHECK NUMBERS
		682	*		
		683	*	PIN ADDRESS ERROR SUBROUTINE	
		684	*		
01154	0 22 4 0004	685	PINERR BPT	4	ERROR STOP PERMITTED
01155	0 0 01 01102	686	BRU	CHECK2	NO
01156	0 0 16 01444	687	LDA	T1	YES
01157	0 0 03 00537	688	BRM	MK0CT	EXPAND ACTUAL TO BCD
01160	0 0 76 00605	689	STA	0TPNM2	SAVE
01161	0 0 74 00606	690	STB	0TPNM2+1	

Table 4-6. 9300 Computer, Extended Mode I/O Test Program (Sheet 16 of 20)

01162	0 0 16 01445	691	LDA	T1+1	EXPAND EXPECTED TO BCD
01163	0 0 03 00537	692	BRM	MK0CT	
01164	0 0 76 00577	693	STA	0TPNM1	SAVE
01165	0 0 74 00630	694	STB	0TPNM1+1	
01166	0 0 21 01241	695	EXU	RDIS	DISCONNECT READER
01167	0 02 0 02641	696	TYP	0,1,4	
01170	1 17 01672	697	LDX	=-15**0177777	
01171	1 30 00574	698	MIB	0TPNM+15,X	OUTPUT GENERAL MESSAGE
01172	1 57 01171	699	BRX	\$-1	
01173	0 30 01624	700	MIB	ERMSG0+4	OUTPUT BLOCK NO.
01174	0 30 01625	701	MIB	ERMSG0+5	
01175	0 30 01626	702	MIB	ERMSG0+6	
01176	0 0 01 00533	703	BRU	0TPIN2	GO OUTPUT RECEIVED AND EXPECTED
		704	*		
		705	*	START TAPE READER SUBROUTINE	
		706	*		
01177	0 0 00 00030	707	STARTP	PZE	
01200	1 17 01660	708	LDX	=-64**0177777	
01201	0 40 37711	709	CLR		
01202	0 1 76 01444	710	STA	BUFFER+64,X	CLEAR BUFFER
01203	1 57 01202	711	BRX	\$-1	
01204	0 0 21 01240	712	EXU	RPT	START TAPE
01205	0 0 41 01177	713	BRR	STARTP	EXIT
		714		PAGE	
		715	*		
		716	*	I/O CHANNEL COMMANDS	
		717	*		
R 01206	002 146 0 30	718	PI0SP	I0SP	IMAGE,64
01207	0100 01244				
R 01210	002 142 0 30	719	PI0SD	I0SD	IMAGE,64
01211	0100 01244				
		720	*		
R 01212	002 142 0 30	721	RI0SD	I0SD	BUFFER,64
01213	0100 01344				
R 01214	002 146 0 30	722	RI0SP	I0SP	BUFFER,65
01215	0101 01344				
R 01216	002 146 0 30	723	RI0SP1	I0SP	BUFFER,32
01217	0040 01344				
R 01220	002 144 0 30	724	RI0RP	I0RP	BUFFER+32,33
01221	0041 01404				
R 01222	002 140 0 30	725	RI0RD	I0RD	BUFFER,56
01223	0070 01344				
		726	*		
		727	*		
		728	*	I/O CHANNEL INSTRUCTIONS.	
		729	*		
01224	0 20 14000	730	PCAT	CAT	0
01225	0 20 12000	731	PCZT	CZT	0

Table 4-6. 9300 Computer, Extended Mode I/O Test Program (Sheet 17 of 20)

01226	0 20 11000	732	PCET	CET	0	
		733	*			
01227	0 02 0 02044	734	PPT	PPT	0,1,1	
01230	0 02 0 00044	735	PTL	PTL	0,1,1	
01231	0 02 00000	736	PDIS	DSC	0	
01232	0 02 12000	737	PASC	ASC	0	
01233	0 02 50000	738	PALC	ALC	0	
		739	*			
		740	*			
01234	0 20 14000	741	RCAT	CAT	0	
01235	0 20 12000	742	RCZT	CZT	0	
01236	0 20 11000	743	RCET	CET	0	
01237	0 20 10400	744	RCIT	CIT	0	
		745	*			
01240	0 02 0 02034	746	RPT	RPT	0,1,1	
01241	0 02 00000	747	RDIS	DSC	0	
01242	0 02 12000	748	RASC	ASC	0	
01243	0 02 50000	749	RALC	ALC	0	
		750	*			
		751	*			
		752	*	OUTPUT IMAGE AREA, INPUT BUFFER AREA		
		753	*			
01244		754	IMAGE	RES	64	
01344		755	BUFFER	RES	64	
		756	*			
		757	*			
		758	*	TEMPORARY STORAGE AND FLAGS		
		759	*			
01444		760	T1	RES	13	
01461		761	SPF	RES	1	IOSP INPUT FLAG
01462		762	PRF	RES	1	PUNCH REPEAT FLAG
		763		PAGE		
		764	*			
		765	*	ERROR AND STATUS MESSAGES.		
		766	*			
01463	52233021	767	DISMSG	BCD		52,1CHANNEL ERROROUSLY DISCONNECTED BEFORE C=0, DURING
01464	45452543					
01465	12255151					
01466	46514664					
01467	62437012					
01470	24316223					
01471	46454525					
01472	23632524					
01473	12222526					
01474	46512512					
01475	23130073					
01476	12246451					
01477	31452712					

Table 4-6. 9300 Computer, Extended Mode I/O Test Program (Sheet 18 of 20)

01500	31466247	768 ERMSG1 BCD	16, IOSP, OUTPUT. 11
01501	73124664		
01502	63476463		
01503	33125252		
01504	31466224	769 ERMSG2 BCD	16, IOSD, OUTPUT. 11
01505	73124664		
01506	63476463		
01507	33125252		
01510	52233021	770 ERMSG4 BCD	52, CHANNEL DID NOT DISCONNECT WHEN C=0 ON IOSD INPUT!
01511	45452543		
01512	12243124		
01513	12454663		
01514	12243162		
01515	23464545		
01516	25236312		
01517	66302545		
01520	12231300		
01521	12464512		
01522	31466224		
01523	12314547		
01524	64635212		
01525	31466224	771 ERMSG5 BCD	12, IOSD, INPUT!
01526	73123145		
01527	47646352		
01530	52233021	772 ERMSG6 BCD	52, CHANNEL DISCONNECTED DURING IOSP INPUT, CIT NEVER
01531	45452543		
01532	12243162		
01533	23464545		
01534	25236325		
01535	24122464		
01536	51314527		
01537	12314662		
01540	47123145		
01541	47646373		
01542	12233163		
01543	12452565		
01544	25511212		
01545	63516425	773	BCD 8, TRUE. 1
01546	33125212		
01547	52246451	774 ERMSG7 BCD	48, DURING IOSP INPUT C=0 INDICATING EOR PAST BUT C
01550	31452712		
01551	31466247		
01552	12314547		
01553	64631223		
01554	13001231		
01555	45243123		
01556	21633145		
01557	27122546		
01560	51124721		

Table 4-6. 9300 Computer, Extended Mode I/O Test Program (Sheet 19 of 20)

01561	62631222			
01562	64631223			
01563	31631266	775	BCD	20. IT WAS NEVER TRUE. I
01564	21621245			
01565	25652551			
01566	12635164			
01567	25331252			
01570	31466247	776	ERMSG8 BCD	12. I'NSP. INPUT!
01571	73123145			
01572	47646352			
01573	52314651	777	ERMSG9 BCD	52. I'IRD ON INPUT DID NOT IGNORE THE LAST 8 CHARACTERS.
01574	24124645			
01575	12314547			
01576	64631224			
01577	31241245			
01600	46631231			
01601	27454651			
01602	25126330			
01603	25124321			
01604	62631210			
01605	12233021			
01606	51212363			
01607	25516273			
01610	12226463	778	BCD	24. BUT READ THE FOLLOWING!
01611	12512521			
01612	24126330			
01613	25122646			
01614	43434666			
01615	31452752			
01616	52524721	779	ERMSGP BCD	8. I'IPARITY
01617	51316370			
01620	52314547	780	ERMSG0 BCD	40. I'INPUT ERROR IN BLOCK NO. NIEXPECTED DD!
01621	64631225			
01622	51514651			
01623	12314512			
01624	22434623			
01625	42124546			
01626	33124552			
01627	25674725			
01630	23632524			
01631	12242452			
01632	51252325	781	BCD	12. RECEIVED DD!
01633	31652524			
01634	12242452			
01635	0 30 00000	782 *		
		783 MIBX	MIB	0
		784 *		
		785 *		
				OUTPUT INSTRUCTION

Table 4-6. 9300 Computer, Extended Mode I/O Test Program (Sheet 20 of 20)

		786	END	BEGIN
01636	00000200			
01636	00000077			
01637	00606047			
01640	00606051			
01641	00606062			
01642	00000070			
01643	00177774			
01644	50277677			
01645	00177775			
01646	57737677			
01647	50277777			
01650	77777776			
01651	00002000			
01652	00177766			
01653	70277777			
01654	00000004			
01655	00000003			
01656	00040000			
01657	00000001			
01660	00177700			
01661	01000000			
01662	00001500			
01663	52121225			
01664	51514651			
01665	12246451			
01666	31452712			
01667	00001504			
01670	77777777			
01671	00001344			
01672	00177761			
01673	00177764			
01674	00177770			
01675	00000000			
01676	00177763			
01677	00001525			
01700	33120152			
01701	00001444			
01702	00177757			
01703	33120252			
01704	00001570			
01705	33120352			
01706	33120452			
01707	00001434			
01710	00000707			
01711	52120000			
01712	00177755			
01713	12000052			
01714	00177767			

SECTION V TROUBLESHOOTING

5.1 GENERAL

5.2 This section contains information useful when troubleshooting the Model 932XX series TMCCs.

5.3 Troubleshooting information contained herein is based on the test programs given in Section 4 of this manual. When an error is indicated during performance of the tests given in Section 4, reference should first be made to the Programming Flow Charts illustrated in figure 5-1 and then to the applicable referenced data.

5.4 TEST PROGRAM FLOW CHART

5.5 Figure 5-1 illustrates the programming flow data for the extended mode test programs given in Section 4. An example of the usage of the flow charts is given in the following paragraphs.

5.6 FLOW CHART EXAMPLE

5.7 In presenting the example of usage of the flow charts, the following points will be assumed:

- a. Paper tape reader connected to one of the interlaced channels.
- b. Breakpoint 1 switch reset.
- c. Breakpoint 2 switch reset.
- d. Breakpoint 3 switch set.
- e. Breakpoint 4 switch reset.
- f. An error exists in block 4.

5.8 At the initialization of the test, the channel number has been inserted on the typewriter, the unit number being used, and the character "R" for the paper tape reader. The program begins at the top of sheet 1 of figure 5-1. The program initializes restart location, disconnects all channels, addresses the keyboard and reads the character typed. As the character typed is an "R", the flow proceeds to the right from CHAR: R to = R and is picked up again on sheet 2. The program then builds channel and unit mask words, and builds: RPT, RCAT, RCET, RCZT, RCIT, RALC, and RDIS (Read Paper Tape, Channel Active Test, Channel Error Test, Channel Zero Count Test, Channel Inter-Record Test, Alert Channel, and Disconnect Channel, respectively). The EOM/EOD commands are then constructed and the keyboard addressed to determine if the letter "S" has

been inserted to start the test. The program returns to sheet 1, KYBD, where the keyboard is addressed, the character "S" is read and Breakpoint Switch 3 is inter-rogated. As Breakpoint Switch 3 is set (paragraph 5.7d) the line S is followed to IN.

5.9 The program proceeds to sheet 8 (circle labeled IN), the tape is started, IOSD 64 is loaded (block 1), a check is made to determine if the count reaches zero and the channel is inactive. As the count has reached zero and the channel is not active, the program then proceeds to check the data as given on sheet 12. The subroutine is then performed to check End Address, Parity, and Input Data. The block number is then saved, no error exists, the channel address is stored and checked against the expected and the input data is compared. As the input data does compare, the program exits from the subroutine and is picked back up again on sheet 8 and proceeds to IN 1.

5.10 IN 1 continues on sheet 9, the tape is started again and IOSP 65 is loaded (block 2). The program checks that the word count does not reach zero and the inter-record test occurs (CIT). It then proceeds from CIT? to IN 2.

5.11 The program then checks to determine if an error exists (sheet 10) and stops the tape before the data is checked. While the tape is stopped, the data check subroutine is performed (sheet 12). After comparison of the data, block no. 2 is entered and the program exists from the subroutine and proceeds (sheet 10) to start the tape (sheet 13) and then loads IOSP 32 (block 3).

5.12 Subsequent to loading IOSP, the subroutine Wait For Count Zero (sheet 6) is again performed. When the count equals zero, the program exits from the subroutine and returns to the main program (sheet 10). As no error existed and the channel is not active, the tape is started again (sheet 13) and IORP 33 is loaded, SPF is reset, and the inter-record indicator (CIT, sheet 9) is turned on at the end of the record. The count should not reach zero and the channel should remain active. As no error exists (sheet 10), the tape is stopped and SPF is inter-rogated. SPF has been reset and the program then continues to IN 3 (sheet 11). The data is then checked (sheet 12) and as it does compare, the program exits from the subroutine, block no. 3 is entered, and the tape is started again (sheet 11). The program then loads IORD 56 (block 4), waits for the channel to be inactive and checks to determine if the channel ignored the last eight characters. As the last eight words were not ignored (paragraph 5.7f), an error exists and the program

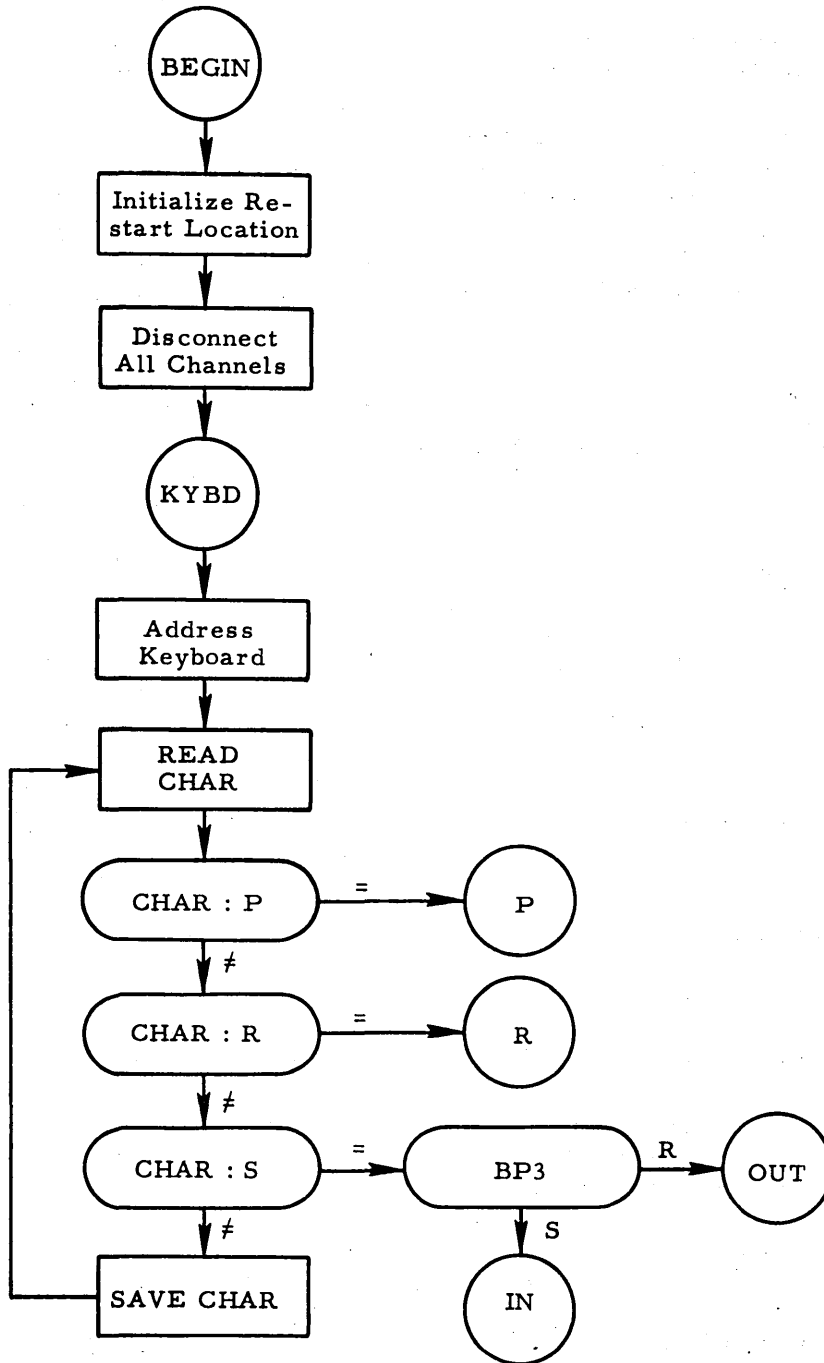


Figure 5-1. Test Program Flow Chart (Sheet 1 of 13)

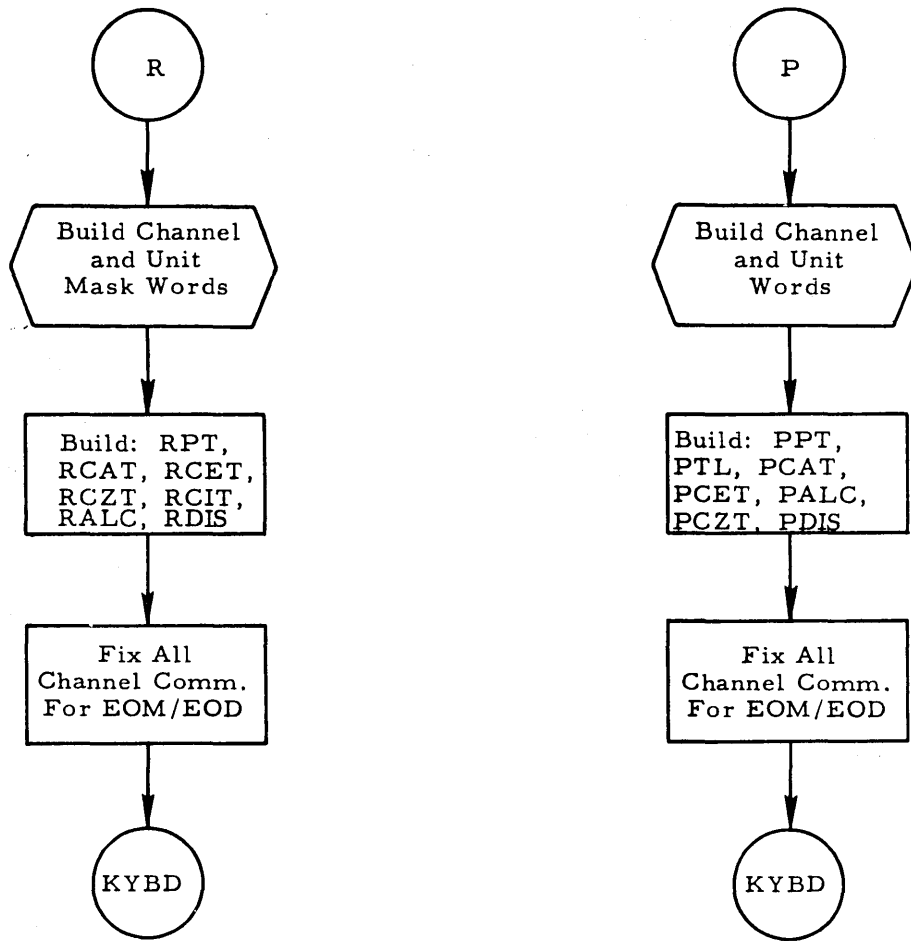


Figure 5-1. Test Program Flow Chart (Sheet 2 of 13)

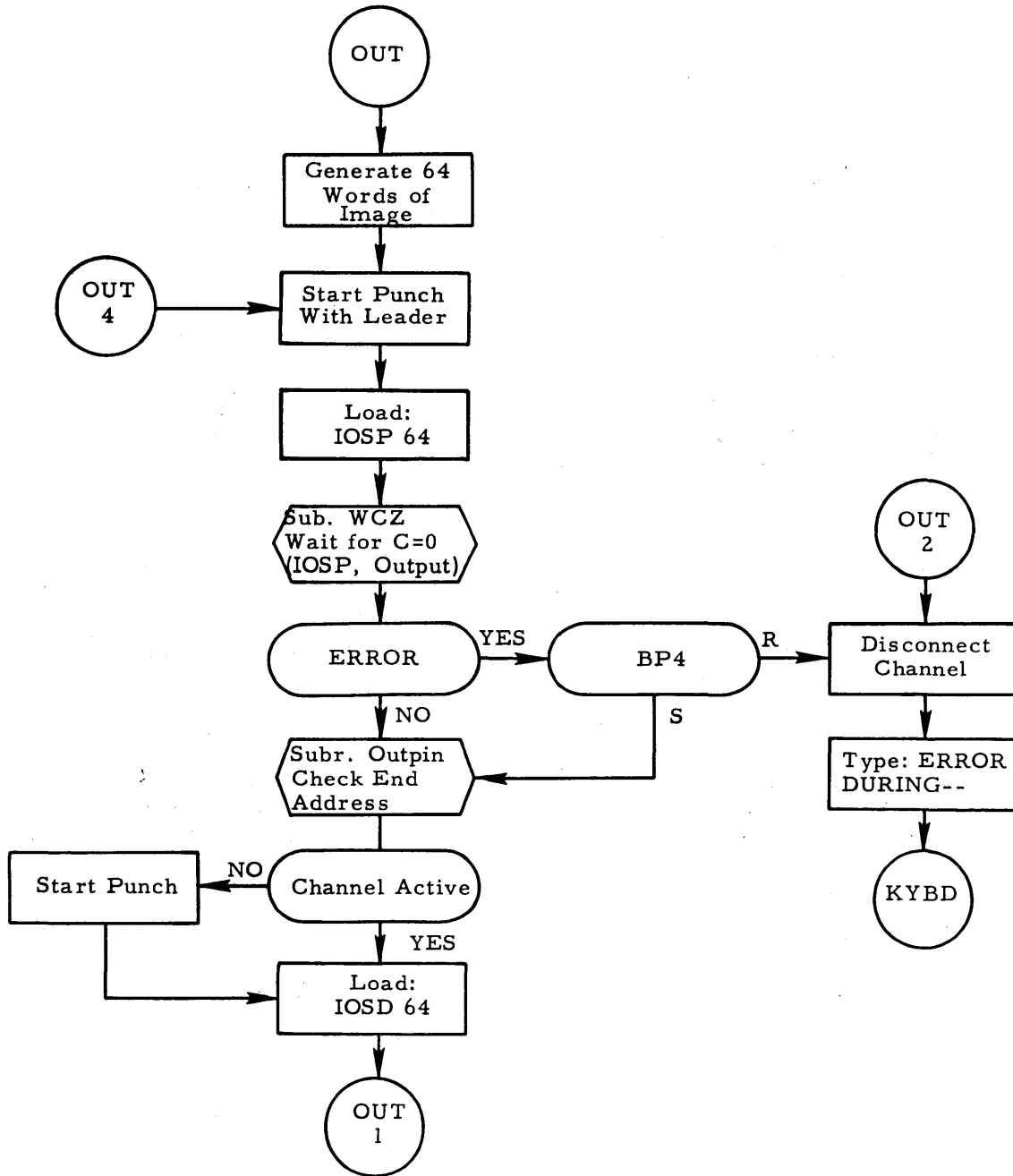


Figure 5-1. Test Program Flow Chart (Sheet 3 of 13)

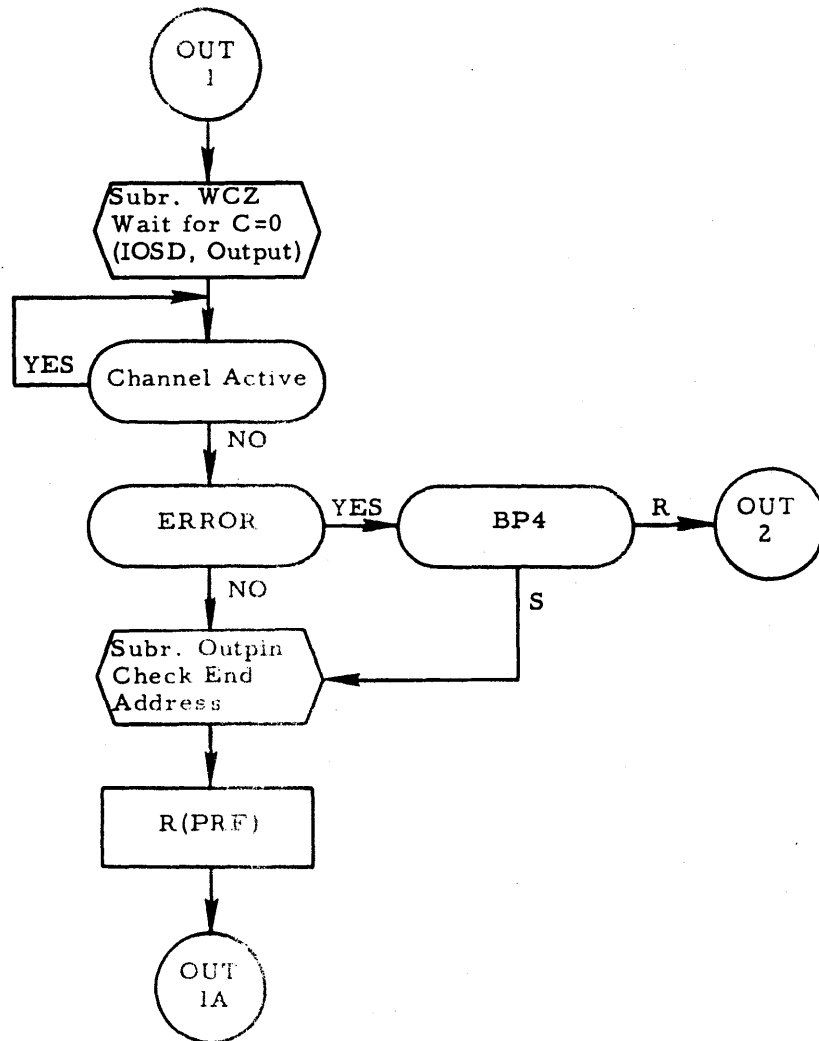


Figure 5-1. Test Program Flow Chart (Sheet 4 of 13)

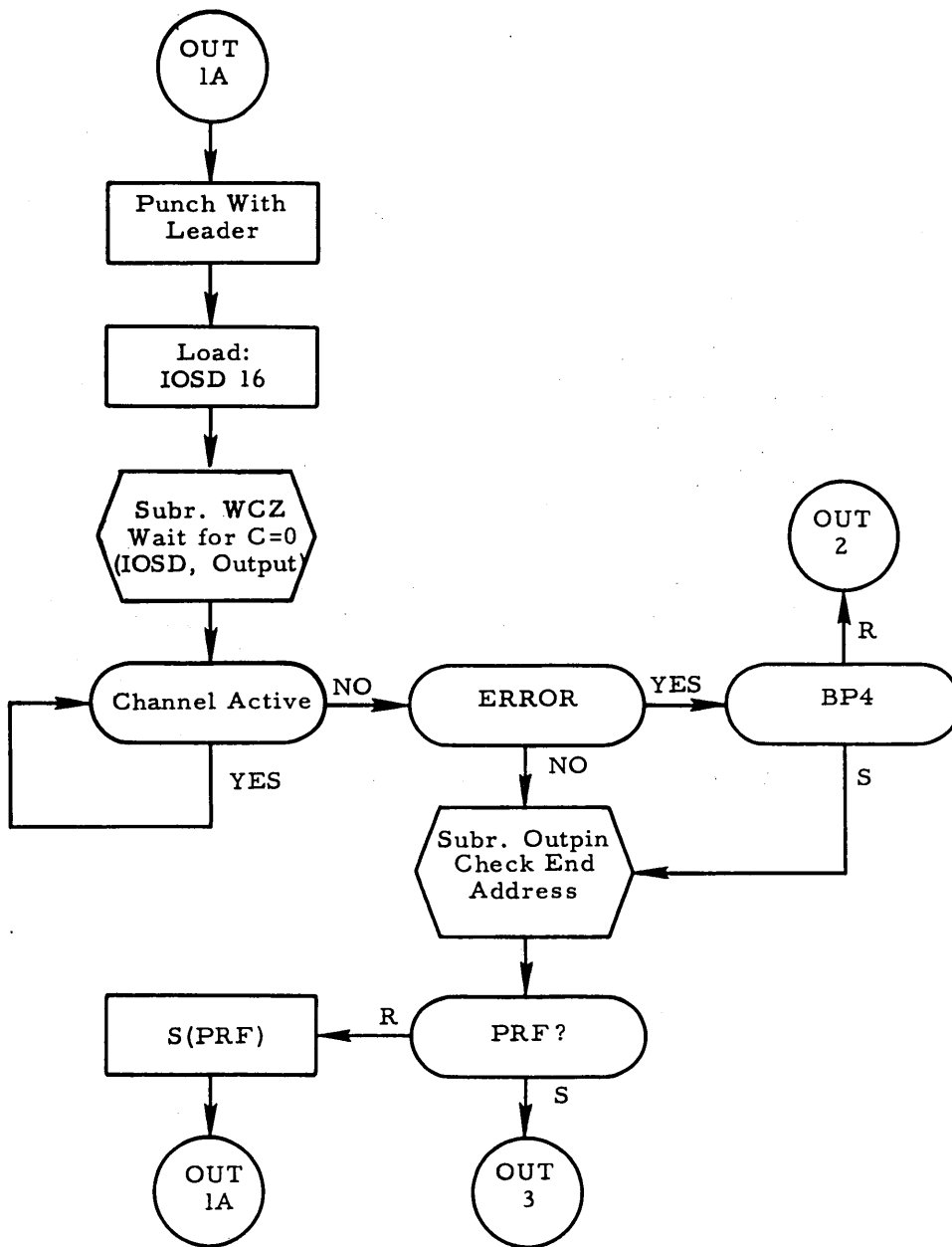


Figure 5-1. Test Program Flow Chart (Sheet 5 of 13)

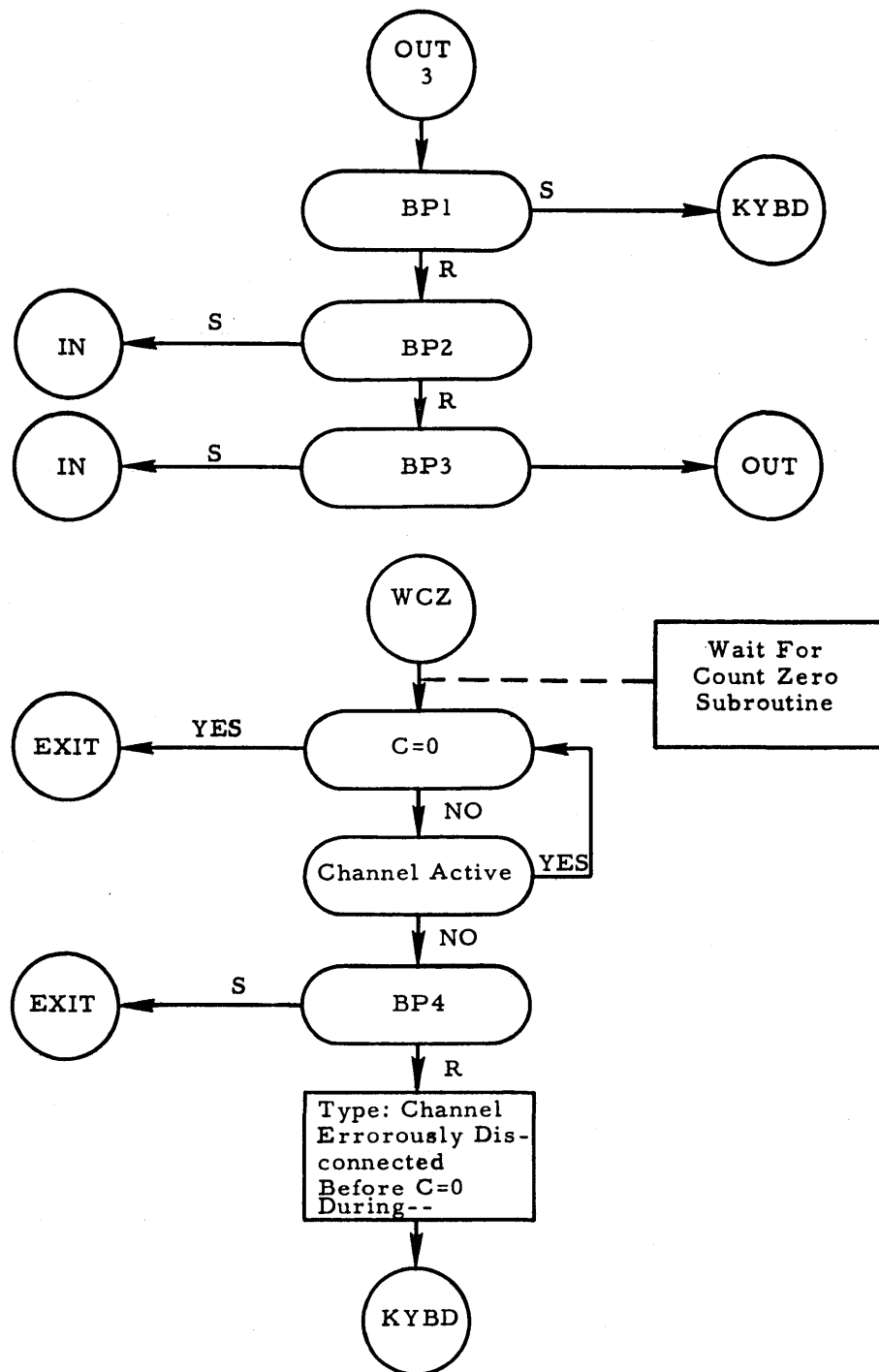


Figure 5-1. Test Program Flow Chart (Sheet 6 of 13)

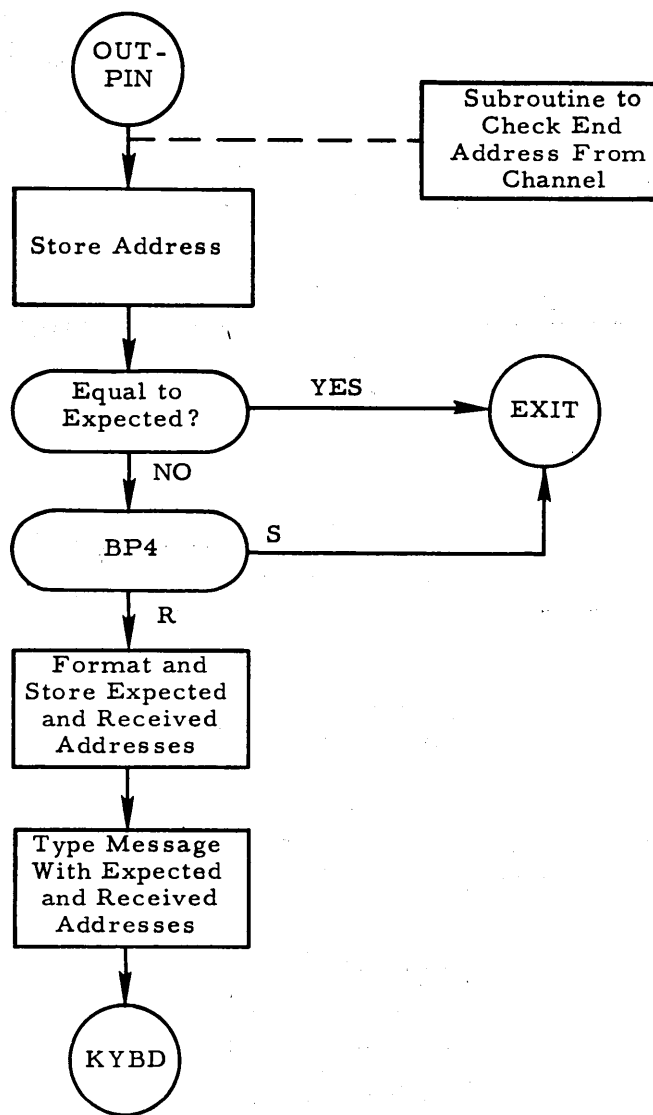


Figure 5-1. Test Program Flow Chart (Sheet 7 of 13)

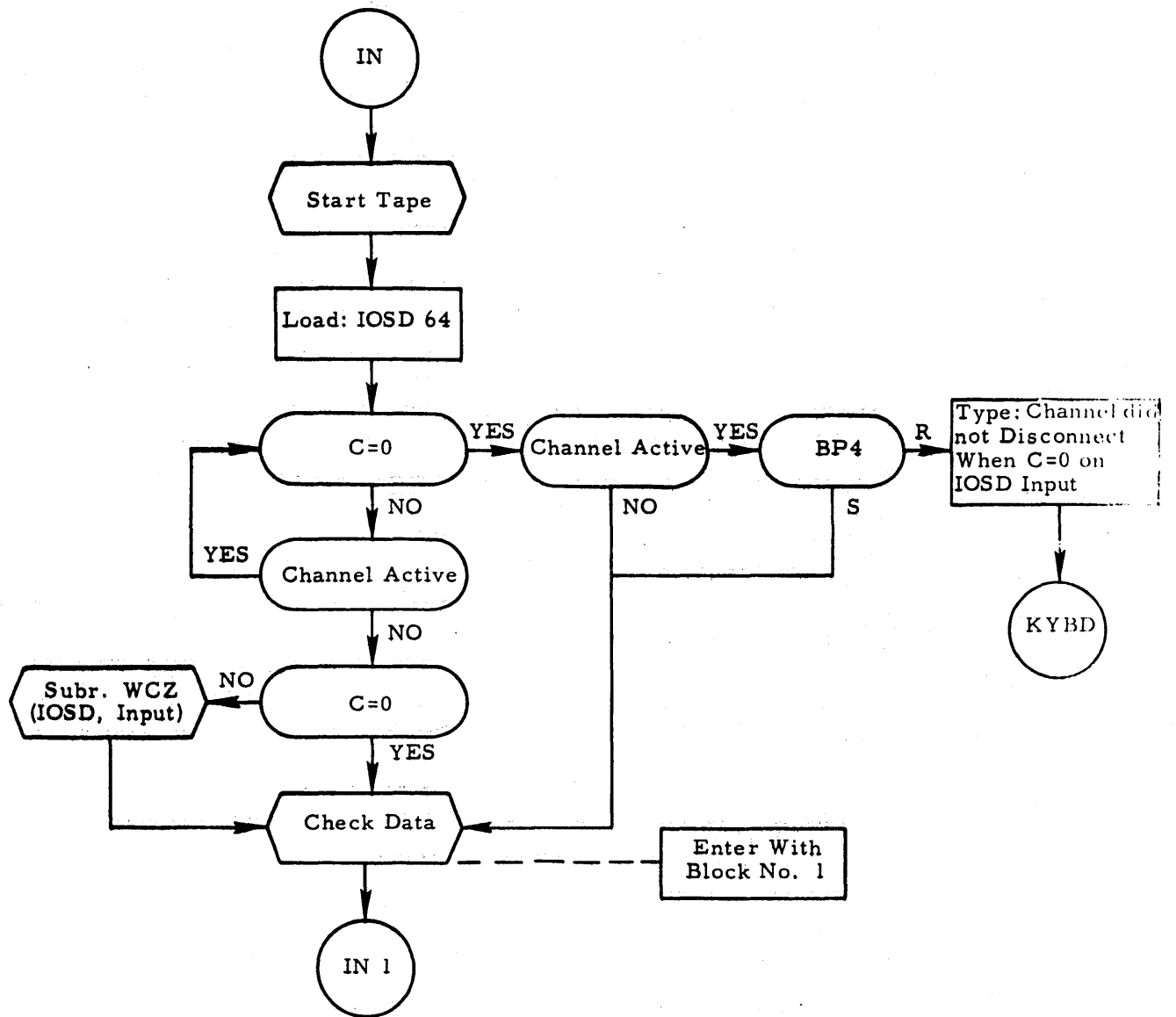


Figure 5-1. Test Program Flow Chart (Sheet 8 of 13)

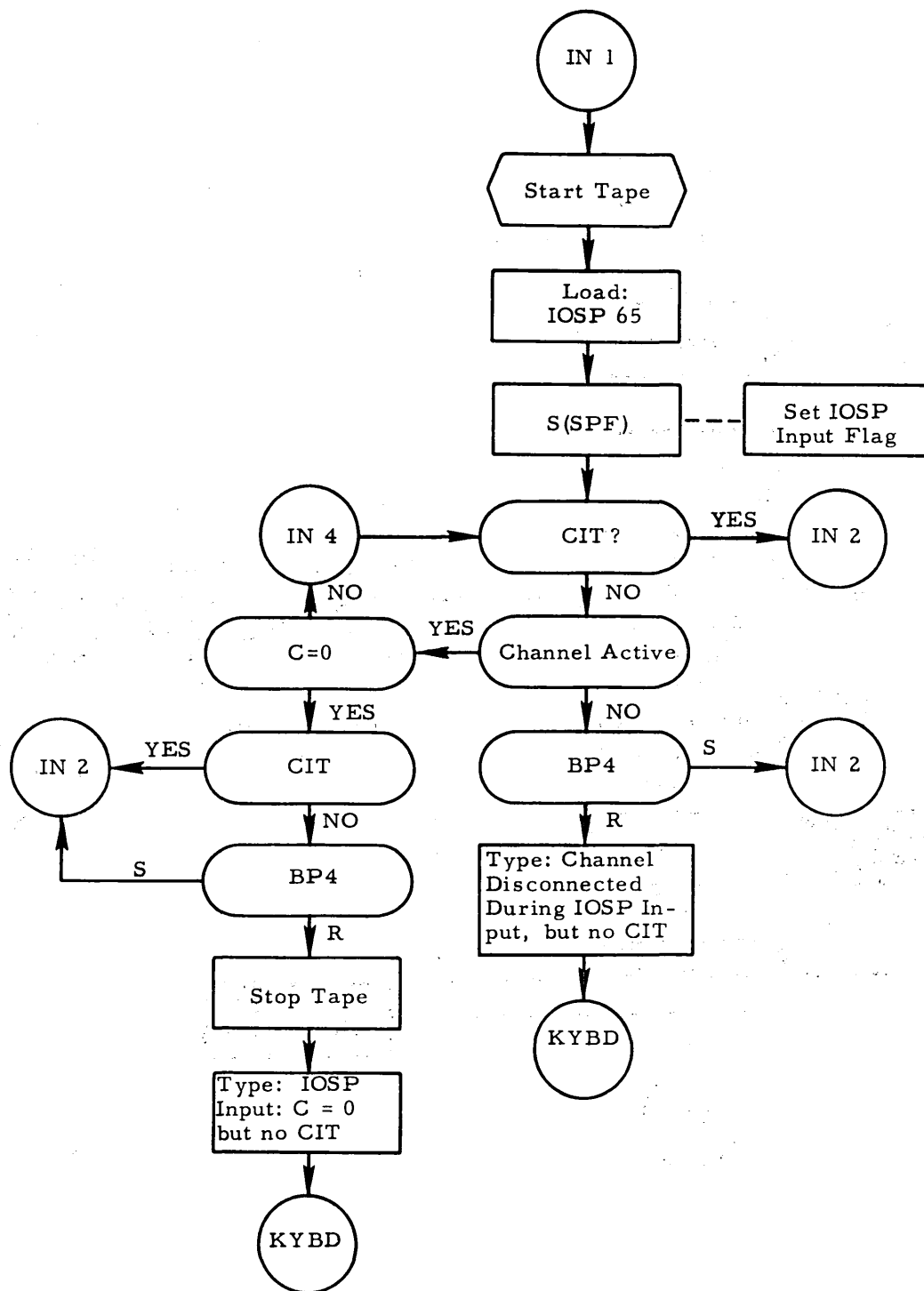


Figure 5-1. Test Program Flow Chart (Sheet 9 of 13)

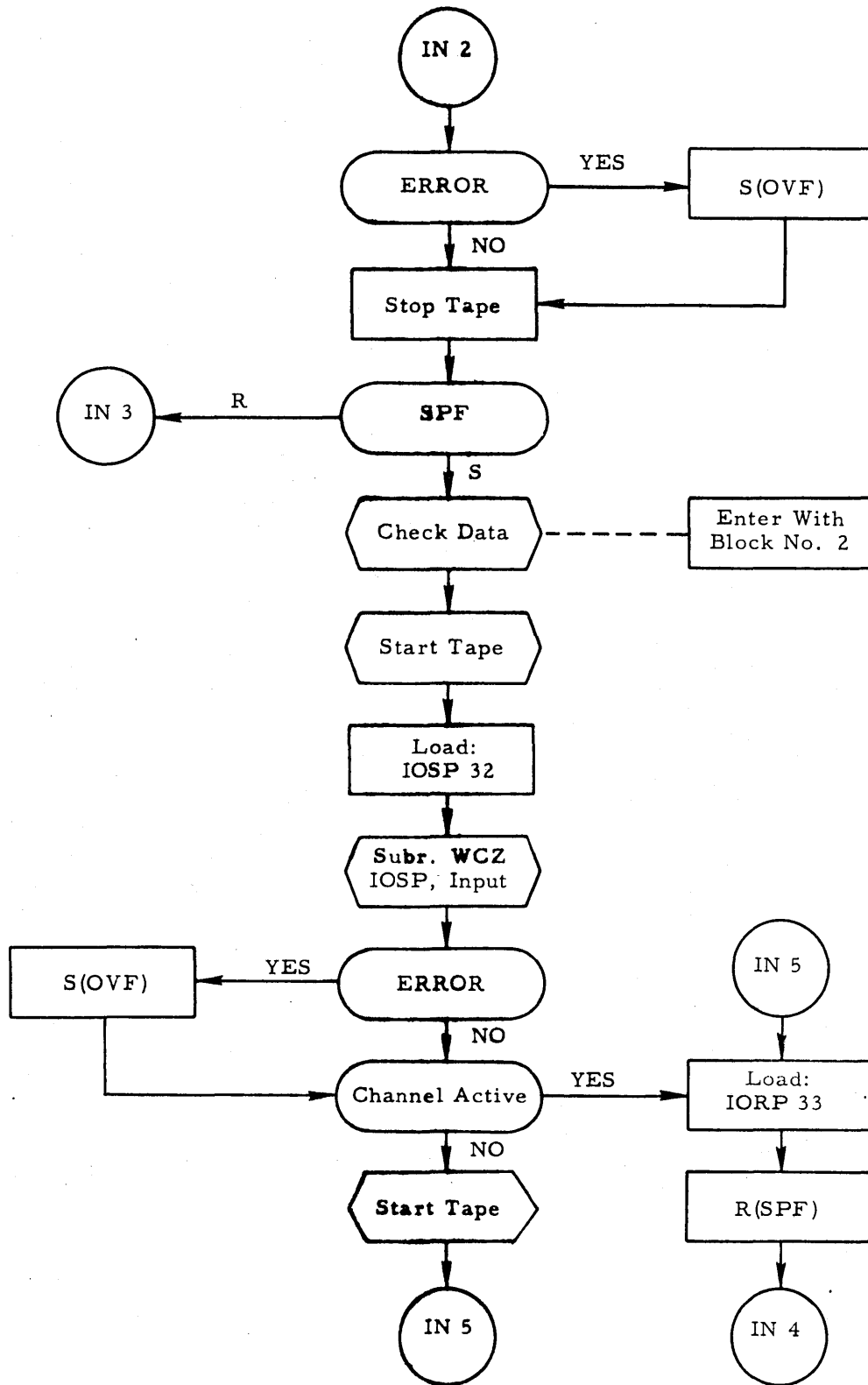


Figure 5-1. Test Program Flow Chart (Sheet 10 of 13)

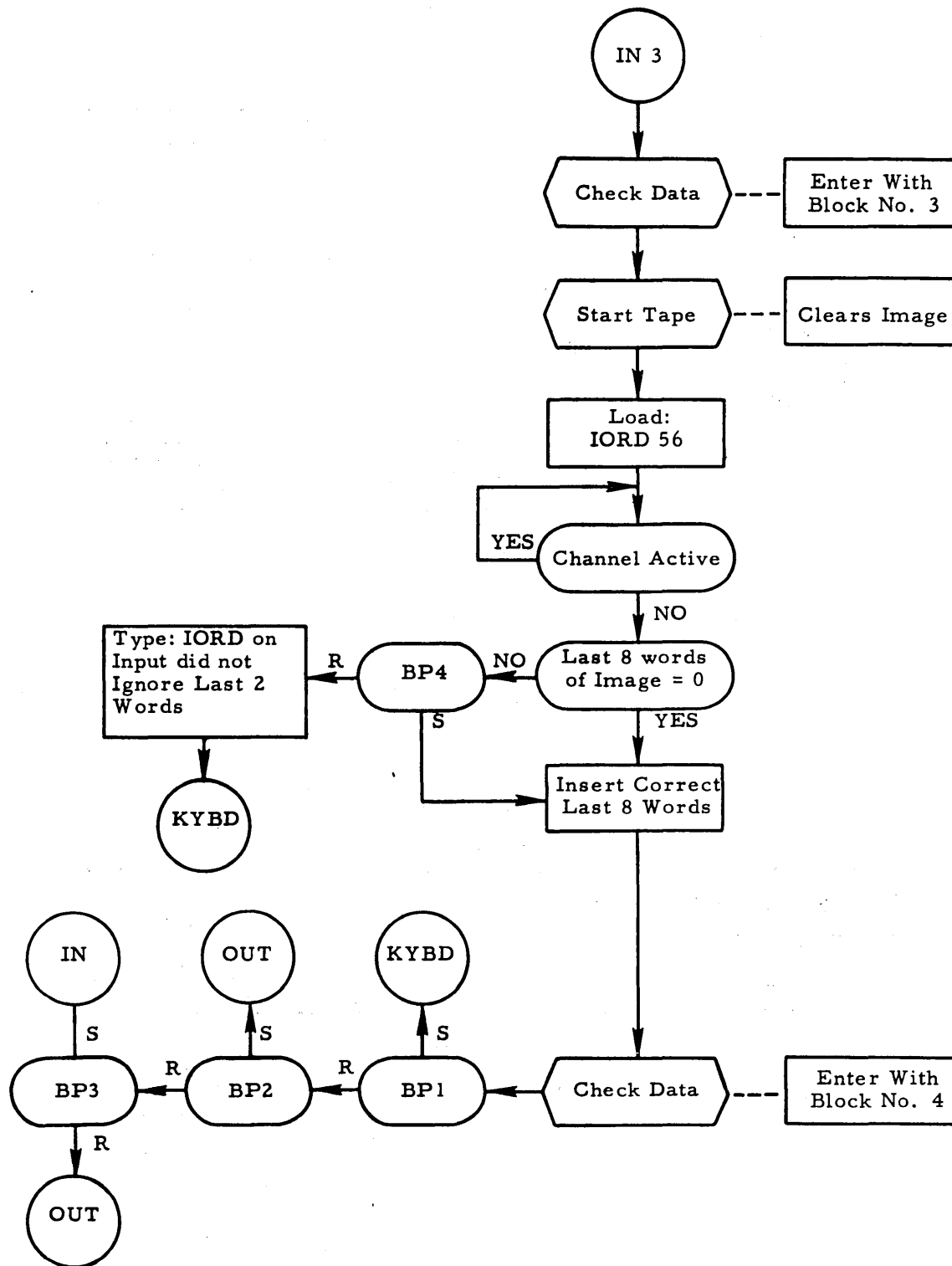


Figure 5-1. Test Program Flow Chart (Sheet 11 of 13)

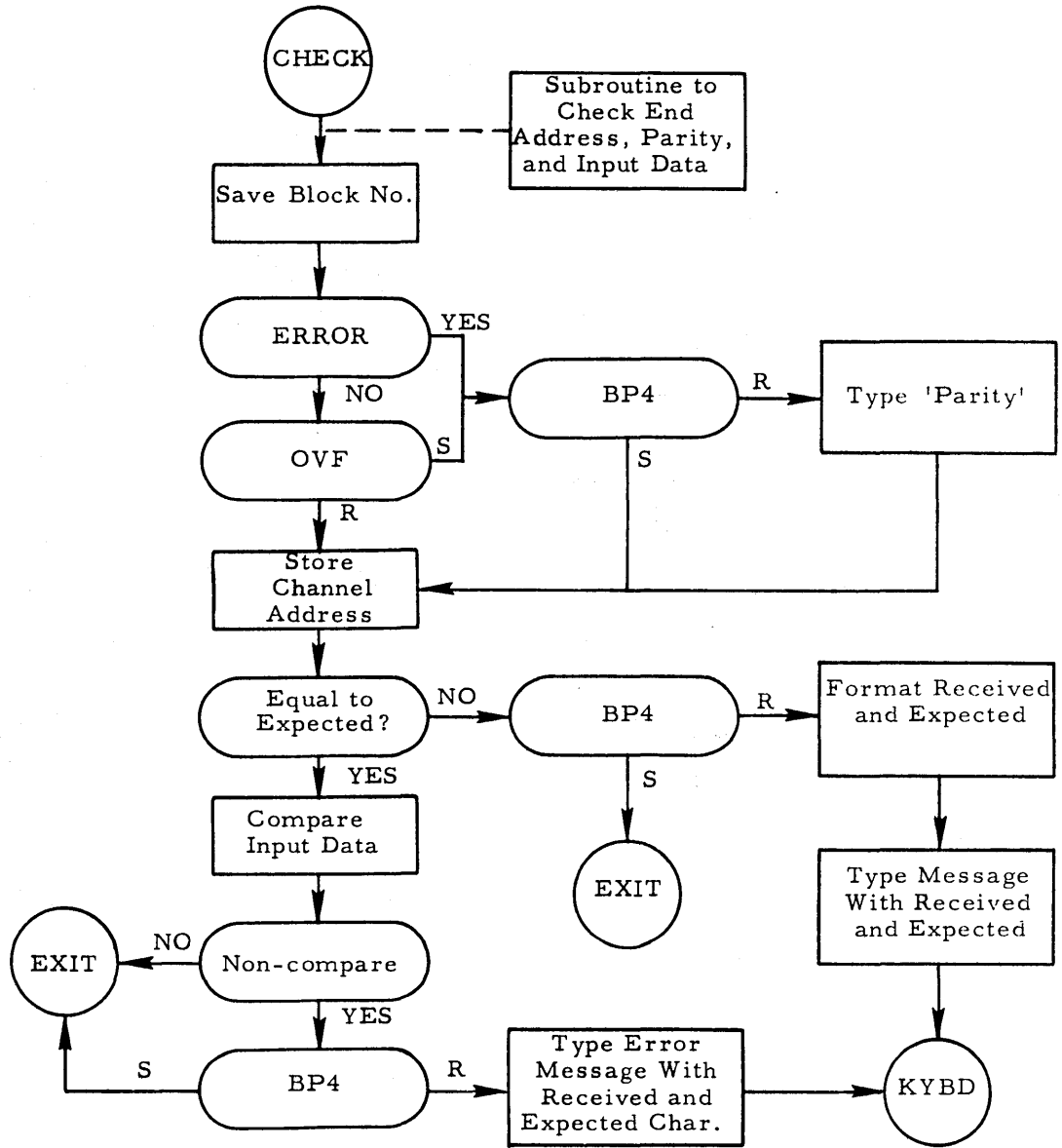


Figure 5-1. Test Program Flow Chart (Sheet 12 of 13)

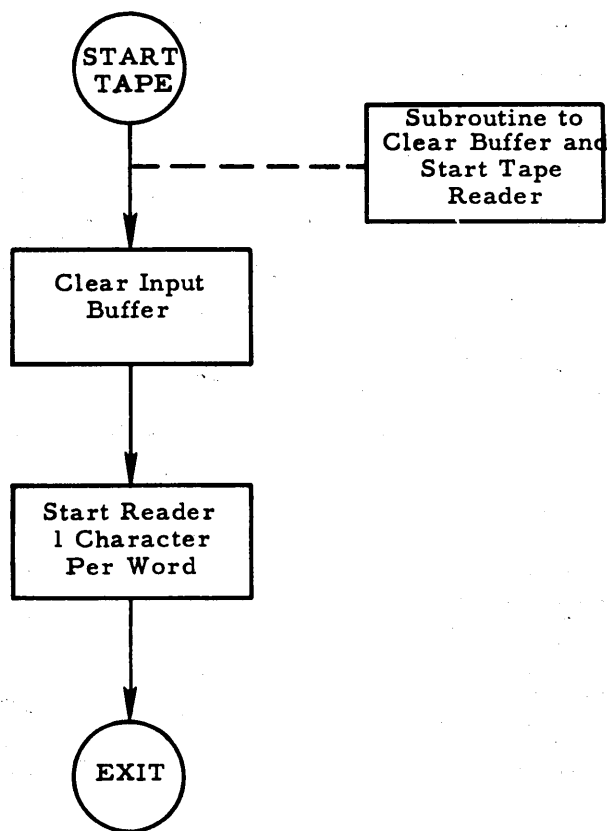


Figure 5-1. Test Program Flow Chart (Sheet 13 of 13)

interrogates Breakpoint Switch 4. Breakpoint Switch 4 is reset (paragraph 5. 7e), therefore, the program types out: IORD on Input did not ignore last two words. Reference should then be made to table 5-8 for information concerning IORD on Input.

5. 13 TROUBLESHOOTING INFORMATION

5. 14 The error printout during the test program is determined by the type of device used with the program. If the device is an output device, the following type of error codes may be printed out:

- a. Typewriter -- Error during IOSD output
- b. Paper Tape -- Error during IORD/IOSD output
- c. Cards -- Error during IOSD/IORD output
- d. Printer -- Error during IOSD/IORD output
- e. Magnetic Tape -- Error during IORD/IORP output.

5. 15 If the device is an input device, the following types of error codes may be printed out:

- a. Typewriter -- Error during IOSD input
- b. Paper Tape -- Error during IORD/IOSD/IORP input
- c. Cards -- Error during IORD/IOSD input
- d. Magnetic Tape -- Error during IORD/IORP input.

5. 16 TROUBLESHOOTING

5. 17 If a failure occurs during performance of the test program and an error message is printed out, determination must be made whether the malfunction is in the central processor unit (CPU), in the input/output device, or in the TMCC. Normally, this can be determined by performing a portion of the applicable test routine for the input/output device and checking for proper operation.

5. 18 If the determination is made that the malfunction is in the TMCC, reference should then be made to the applicable table (tables 5-1 thru 5-10) for that function. The table describes the function and references the paragraphs in the Theory of Operation section (Section 3) where a detailed description of that function is described.

5. 19 The logic equations pertaining to the particular function can be determined from the description in the theory of operations. A comparison of the logic equations and terms will indicate the particular terms peculiar to the function which has failed. Reference to the logic layout drawings (listed in Section 1) will indicate the module in which the term is used, the physical location of the module, and the terminal connections where the term can then be found.

5. 20 Normal troubleshooting procedures can be performed to pinpoint the malfunction to a particular component or terminal.

5. 21 Physical location of components and schematics of each module can be found in Section 6.

Table 5-1. IOSP Output Function, W (A) Channel

Iwg	Iwh	Iwi	Output Function	Sec. 3, Par. Ref.
1	1	1	1. I1w at Iwf if Iwk When the channel counts C down to zero (Iwf), the channel generates a zero word count interrupt (I1w), if armed (Iwk), indicating the program should reload the interlace portion of the channel to continue writing in the same record. Failure to reload the interlace before the buffer transmits all of the characters in its registers and before the peripheral device requests the next character from the buffer sets the channel error indicator.	3. 152 thru 3. 157

Table 5-2. IORP Output Function, W (A) Channel

Iwg	Iwh	Iwi	Output Function	Sec. 3, Par. Ref.
1	1	0	<ol style="list-style-type: none"> 1. I1w at Iwf if Iwk 2. At Iwf, reset W0 3. I2w at Mtgw or Whs W11 if Iwj 4. At Whs W11, disconnect <p>When the channel interlace counts C down to zero (Iwf), the channel generates a zero word count interrupt (I1w), if armed (Iwk), notifying the channel buffer that it has received the last word that is to be output. At zero word count (Iwf), the Halt Interlock flip-flop, W0, is reset to inhibit additional time-share requests. When the device receives the last word from the buffer, it sends an End-of-Record response (Whs W11) back to the buffer. If armed, (Iwj), the buffer generates an End-of-Record Interrupt (I2w) and sets the inter-record indicator. If the device is magnetic tape, an End-of-Record response (tape gap signal, Mtgw,) signal is sent to the buffer but the tape continues to move. If the program does not execute an EOM to write a new tape before the tape gap signal drops, the channel disconnects (Whs W11) and the tape stops.</p>	<p>3. 144 thru 3. 151</p>

Table 5-3. IOSD Output Function, W (A) Channel

Iwg	Iwh	Iwi	Output Function	Sec. 3, Par. Ref.
1	0	1	<ol style="list-style-type: none"> 1. I1w at Iwf if Iwk 2. At Iwf, reset W0 3. Disconnect at Iwf W11 or Whs 4. I2w at disconnect, if Iwj <p>When the channel interlace counts C down to zero (Iwf), the channel generates a zero word count interrupt (I1w), if armed (Iwk), indicating the last character has been transmitted. At zero word count (Iwf), the Halt Interlock flip-flop (W0) is reset inhibiting additional time-share requests. For devices other than magnetic tape (W11), the Halt Detector flip-flop (Wh) is set on reaching zero word count (Iwf) when the last character has been clocked from the buffer. The Halt Detector also sets on occurrence of a Halt Signal (Whs). Setting of the Halt Detector initiates a buffer disconnect sequence. The Signal Complete flip-flop (Wcs) is set and if the End-of-Record Interrupt Enable (Iwj) has been previously armed, an End-of-Record Interrupt (I2w) is generated.</p>	<p>3. 140 thru 3. 143</p>

Table 5-4. IORD Output Function, W (A) Channel

Iwg	Iwh	Iwi	Output Function	Sec. 3, Par. Ref.
1	0	0	<ol style="list-style-type: none"> 1. I1w at Iwf if Iwk 2. At Iwf, reset W0 3. Disconnect at Whs 4. I2w at disconnect, if Iwj <p>When the channel counts C down to zero (Iwf), the channel generates a zero word count interrupt (I1w), if armed (Iwk), indicating that the last characters have been transmitted. At zero word count (Iwf), the Halt Interlock flip-flop (W0) is reset inhibiting additional time-share requests. If Halt Signal (Whs) is received, the Halt Detector (Wh) is set and a disconnect occurs. The Signal Complete flip-flop (Wsc) is set and, if armed (Iwj), an End-of-Record Interrupt (I2w) is generated.</p>	3. 134 thru 3. 139

Table 5-5. IO SP Input Function, W (A) Channel

Iwg	Iwh	Iwi	Input Function	Sec. 3, Par. Ref.
1	1	1	<ol style="list-style-type: none"> 1. I1w at Iwf if Iwk 2. At Mtgw or Whs W11, flush and store last character(s) if Iwf 3. I2w at Mtgw or Whs W11 if Iwj 4. Disconnect at Whs W11 <p>When the channel counts C down to zero (Iwf), the channel generates a zero word count interrupt (I1w), if armed (Iwk), indicating the program should reload the interlace portion of the channel to continue reading the record. If the End-of-Record (Mtgw or Whs W11) occurs before zero word count (Iwf), the buffer is flushed and the completed word is stored in memory. If the End-of-Record Interrupt Enable (Iwi) has been armed, an End-of-Record interrupt (I2w) is generated when a tape gap (Mtgw) or halt signal (Whs) is detected from the device. For magnetic tape operation (W11), a new EOM may be given within one millisecond from the occurrence of I2w to permit the tape system to proceed to a new record. Failure to give an EOM results in the tape stopping and the buffer disconnecting.</p>	3. 176 thru 3. 179

Table 5-6. IORP Input Function, W (A) Channel

Iwg	Iwh	Iwi	Input Function	Sec. 3, Par. Ref.
1	1	0	<ol style="list-style-type: none"> 1. I1w at Iwf if Iwk 2. Inhibit rate errors if Iwf 3. At Mtgw or Whs W11, flush and store last character(s) if Iwf 4. I2w at Mtgw or Whs W11 if Iwj 5. Disconnect at Whs W11 <p>When the channel counts C down to zero (Iwf), the channel generates a zero word count interrupt (I1w), if armed (Iwk), indicating the program should reload the interlace portion of the channel to continue reading the record. Additional characters entering the channel after zero word count are precessed into the W register. Parity and rate errors cannot occur after zero word count because of Iwf. Detection of magnetic tape gap (Mtgw) or a halt signal (Whs) sets the End-of-Record detector. If the End-of-Record detector is set before zero word count has occurred (Iwf), the buffer is flushed and the completed word is stored in memory. If the End-of-Record Interrupt Enable (Iwj) has been armed, an End-of-Record interrupt (I2w) occurs. Failure to reload the interlace within one millisecond of I2w results in the tape stopping and the buffer disconnecting.</p>	3. 168 thru 3. 175

Table 5-7. IOSD Input Function, W (A) Channel

Iwg	Iwh	Iwi	Input Function	Sec. 3, Par. Ref.
1	0	1	<ol style="list-style-type: none"> 1. I1w at Iwf if Iwk 2. At Whs, flush and store last character(s) if Iwf 3. Disconnect at Iwf W11 or Whs 4. I2w at disconnect if Iwj <p>When the channel counts C down to zero (Iwf), the channel generates a zero word count interrupt (I1w), if armed (Iwk). Should an End-of-Record (Whs) occur before zero word count (Iwf) is established, any characters remaining in the W register are flushed and stored. The Halt Detector is now permitted to set by Iwf W11 or Whs W11 and the channel is disconnected. If the End-of-Record Interrupt Enable (Iwj) has been armed, an End-of-Record interrupt (I2w) occurs.</p>	3. 164 thru 3. 167

Table 5-8. IORD Input Function, W (A) Channel

Iwg	Iwh	Iwi	Input Function	Sec. 3, Par. Ref.
1	0	0	<ol style="list-style-type: none"> 1. I1w at Iwf if Iwk 2. Inhibit rate errors if Iwf 3. At Mtgw or Whs W11, flush and store last character(s) if Iwf 4. Disconnect at Whs 5. I2w at disconnect if Iwj <p>When the channel counts C down to zero (Iwf), the channel generates a zero word count interrupt (I1w), if armed (Iwk). Parity and rate errors (We) are inhibited by Iwf after zero word count is established. Should an End-of-Record (Mtgw or Whs W11) occur before zero word count is established (Iwf), the End-of-Record detector is set and any characters remaining in the W register are flushed and stored. The Halt Detector (Wh) is permitted to set by Whs, the buffer is cleared and the channel disconnected. If the End-of-Record Interrupt Enable (Iwj) has been armed, an End-of-Record interrupt (2w) occurs.</p>	<p>3. 158 thru 3. 163</p>

Table 5-9. Output Functions, Y Channel

Iyg	Iyh	Iyi	Output Function	Sec. 3, Par. Ref.
1	1	1	<p>IOSP</p> <ol style="list-style-type: none"> 1. I1y at Iyf if Iyk 	<p>3. 152 thru 3. 157</p>
1	1	0	<p>IORP</p> <ol style="list-style-type: none"> 1. I1y at Iyf if Iyk 2. At Iyf, reset Y0 3. I2y at Mtgy or Yhs Y11 if Iyj 4. At Yhs Y11, disconnect 	<p>3. 144 thru 3. 151</p>
1	0	1	<p>IOSD</p> <ol style="list-style-type: none"> 1. I1y at Iyf if Iyk 2. At Iyf, reset Y0 3. Disconnect at Iyf Y11 or Yhs 4. I2y at disconnect, if Iyj 	<p>3. 140 thru 3. 143</p>
1	0	0	<p>IORD</p> <ol style="list-style-type: none"> 1. I1y at Iyf if Iyk 2. At Iyf, reset Y0 3. Disconnect at Yhs 4. I2y at disconnect, if Iyj 	<p>3. 134 thru 3. 139</p>
<p>The output functions for the Y channel are identical to those of the W channel given in tables 5-1, 5-2, 5-3, and 5-4. The only difference is the substitution of the letter "y" for "w" in the logic terms.</p>				

Table 5-10. Input Functions, Y Channel

Iyg	Iyh	Iyi	Input Functions	Sec. 3, Par. Ref.
1	1	1	IOSP 1. Ily at Iyf if Iyk 2. At Mtgy or Yhs Y11, flush and store last character(s) if Iyf 3. I2y at Mtgy or Yhs Y11 if Iyj 4. Disconnect at Yhs Y11	3. 176 thru 3. 179
1	1	0	IORP 1. Ily at Iyf if Iyk 2. Inhibit rate errors if Iyf 3. At Mtgy or Yhs Y11, flush and store last character(s) if Iyf 4. I2y at Mtgy or Yhs Y11 if Iyj 5. Disconnect at Yhs Y11	3. 168 thru 3. 175
1	0	1	IOSD 1. Ily at Iyf if Iyk 2. At Yhs, flush and store last character(s) if Iyf 3. Disconnect at Iyf Y11 or Yhs 4. I2y at disconnect if Iyj	3. 164 thru 3. 167
1	0	0	IORD 1. Ily at Iyf if Iyk 2. Inhibit rate errors if Iyf 3. At Mtgy or Yhs Y11, flush and store last character(s) if Iyf 4. Disconnect at Yhs 5. I2y at disconnect if Iyj	3. 158 thru 3. 163
The input functions for the Y channel are identical to those of the W channel given in tables 5-5, 5-6, 5-7, and 5-8. The only difference is the substitution of the letter "y" for "w" in the logic terms.				

SECTION VI
DRAWINGS

6.1 GENERAL

6.2 This section contains drawings useful when troubleshooting and maintaining the TMCC.

6.3 SCOPE OF SECTION

6.4 Included in this section are assembly drawings, schematic diagrams, and material lists for each module.

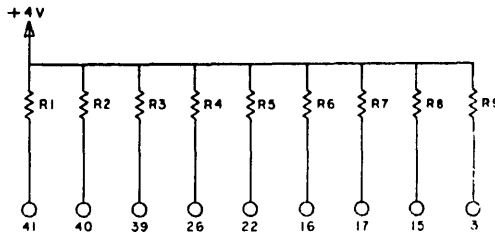
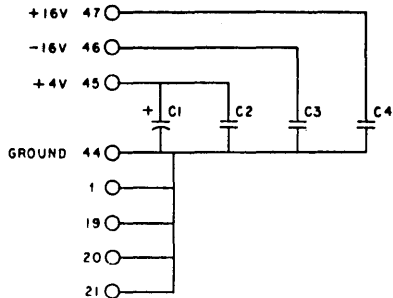
6.5 The type and quantity of each module are listed in Section 1, table 1-4. Physical location of each module is illustrated in Section 4, figure 4-4.

6.6 Also included in this section is the Semiconductor Cross Reference which provides a cross-reference between Scientific Data Systems semiconductor numbers, commercial Electronic Industries Association (EIA) numbers, specification numbers, and replacements for obsolete semiconductors.

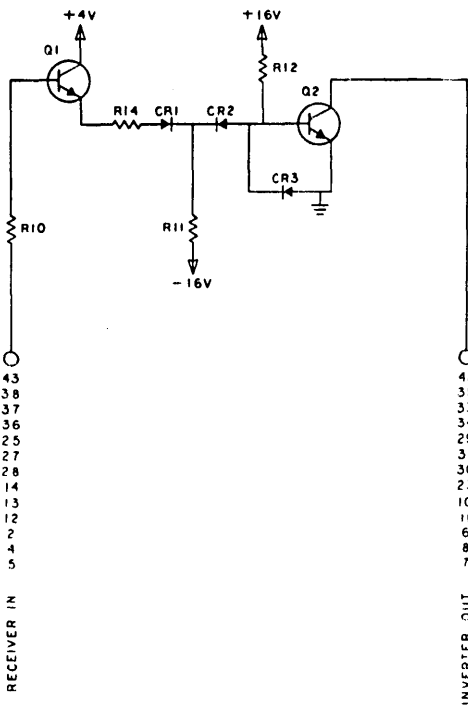
REVISIONS
 103655 A
 RELEASED TO MRS

BRANDING NO.	DESCRIPTION	REFERENCE DESIGNATION
NOTES UNLESS SPECIFIED: 1. ALL DIMENSIONS IN INCHES 2. DIMENSIONS IN PARENTHESES ARE PERMISSIBLE VARIATIONS 3. DIMENSIONS ARE TO CENTER UNLESS OTHERWISE SPECIFIED 4. ALL SURFACES UNLESS OTHERWISE SPECIFIED	SCHEMATIC RECEIVER INVERTER	103655 D
AR 32	AR 32	A

POLARIZING PINS: 18 & 22

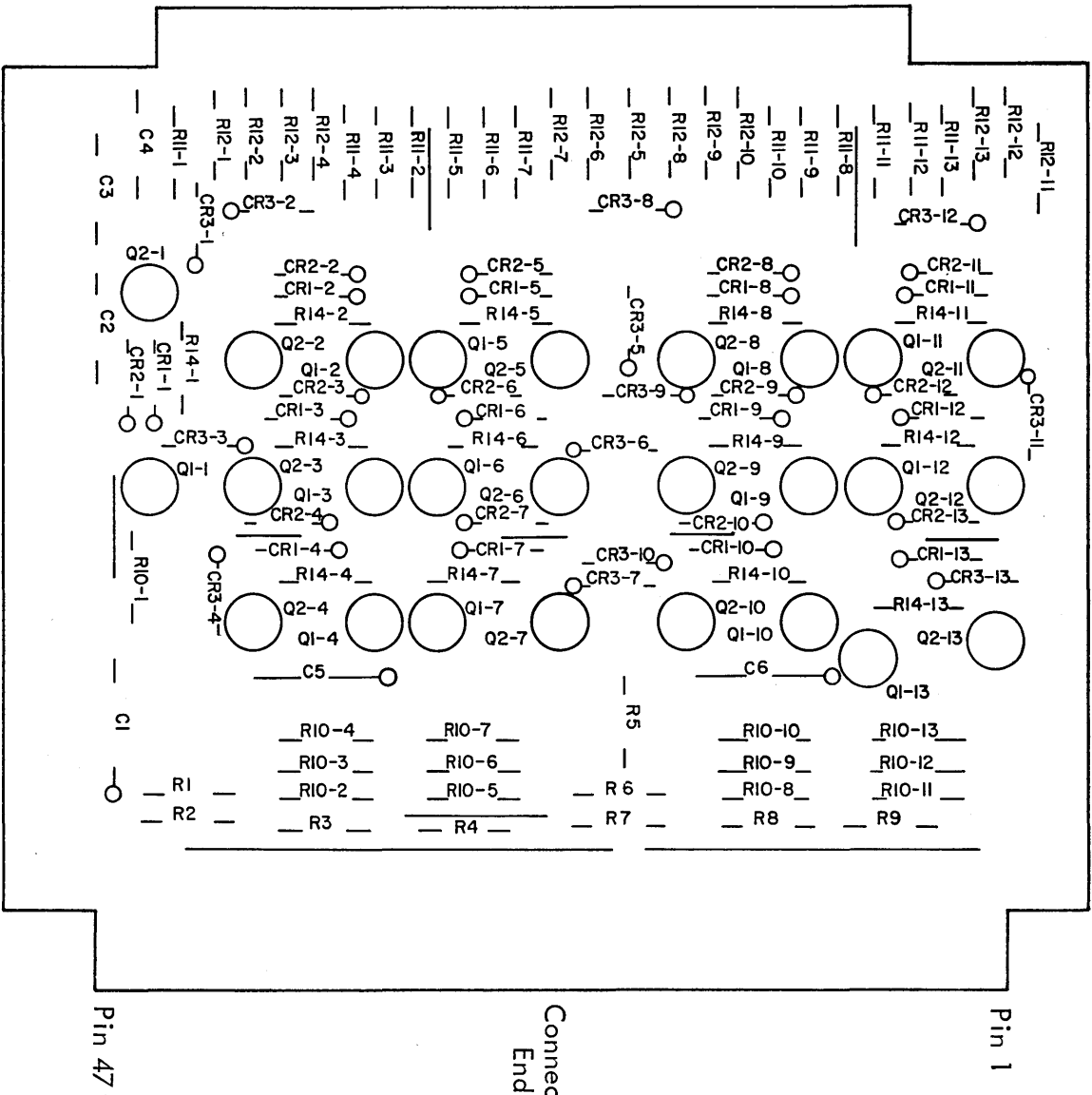


- CIRCUITS
- 1 43
 - 2 38
 - 3 37
 - 4 36
 - 5 25
 - 6 27
 - 7 28
 - 8 14
 - 9 13
 - 10 12
 - 11 2
 - 12 4
 - 13 5



RECEIVER IN

INVERTER OUT

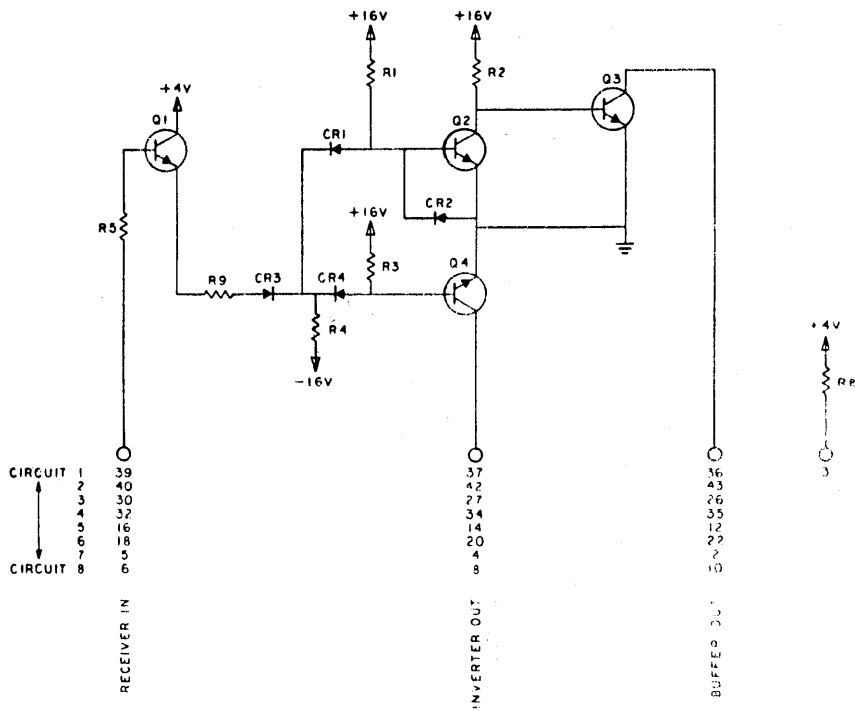
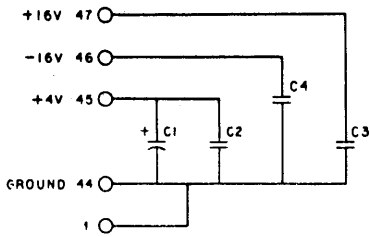


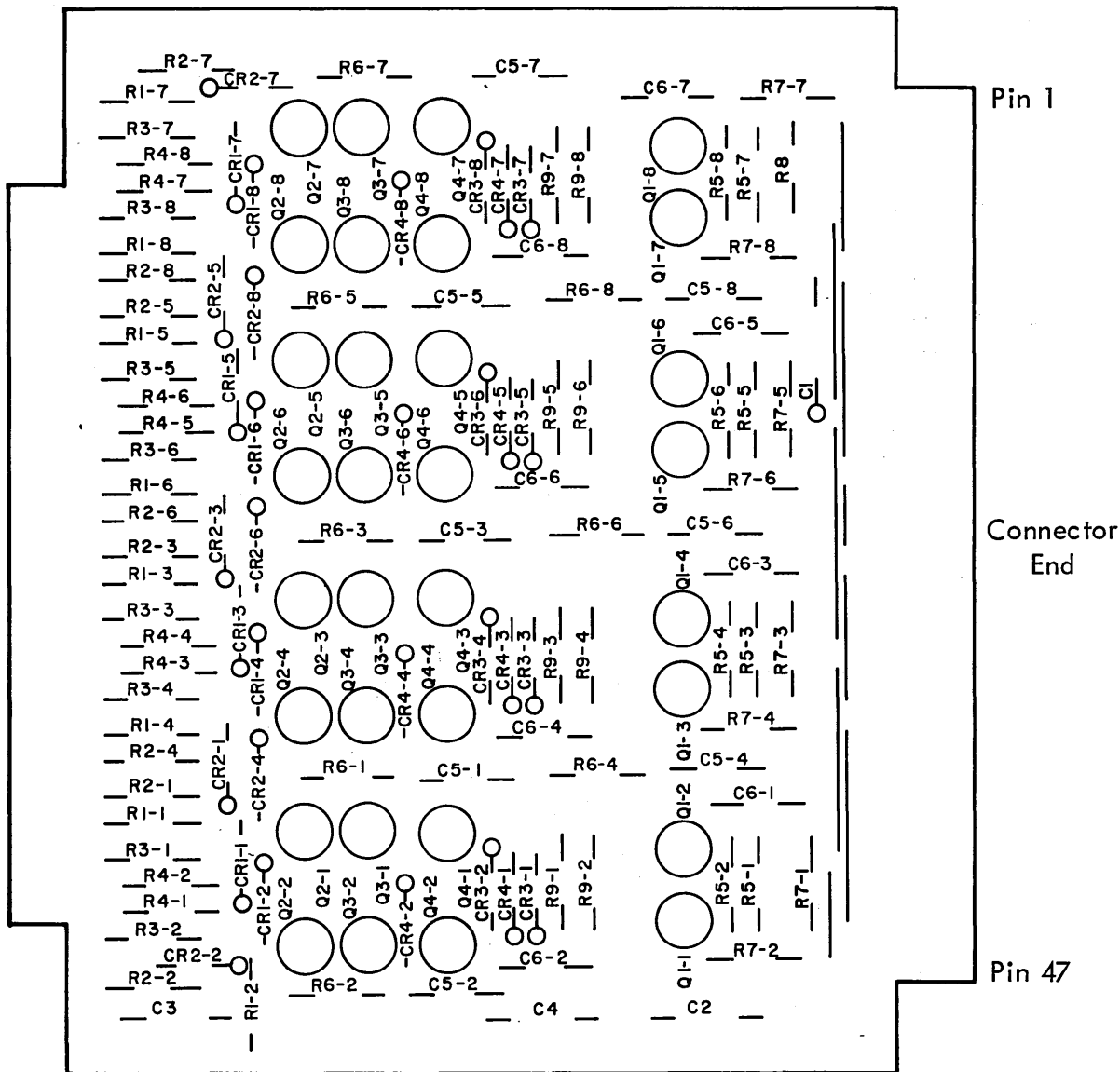
Receiver Inverter AB52
103657C

RELEASSED TO THE PUBLIC

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TITLE		RECEIVER INVERTER BUFFER		A	
DRAWN BY		CHECKED BY		DATE	
DESIGNED BY		APPROVED BY		SCALE	
PART NO.		REV.		REV.	
A853		1		10/3/52	

POLARIZING PINS: 20 & 44



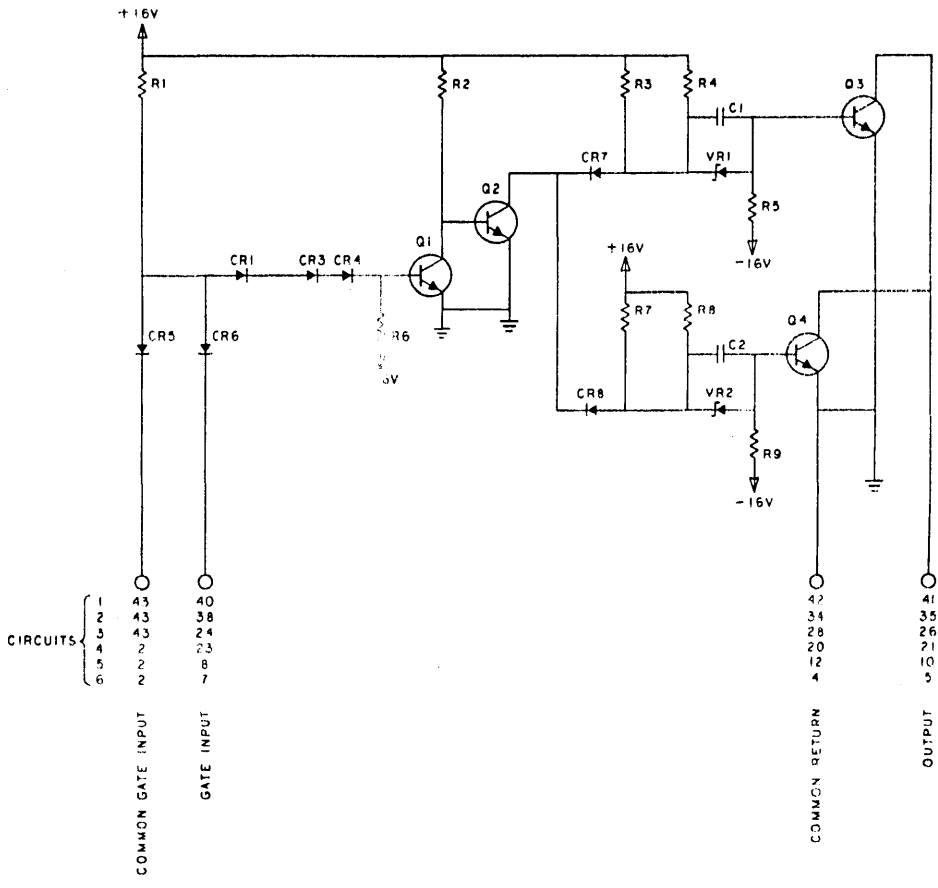
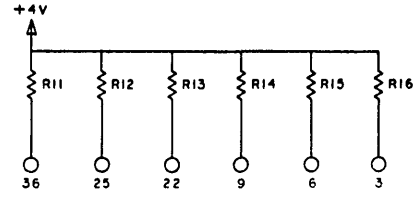
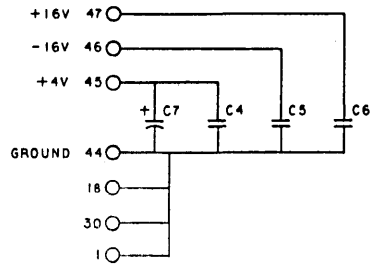


Receiver Inverter Buffer AB53
103726A

REVISIONS	
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A	RELEASED TO MFG
B	SEE REVISION E.O.
C	SEE REV E.O.

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APPROVED BY: [Signature]		MATERIAL CITY		SCALE	

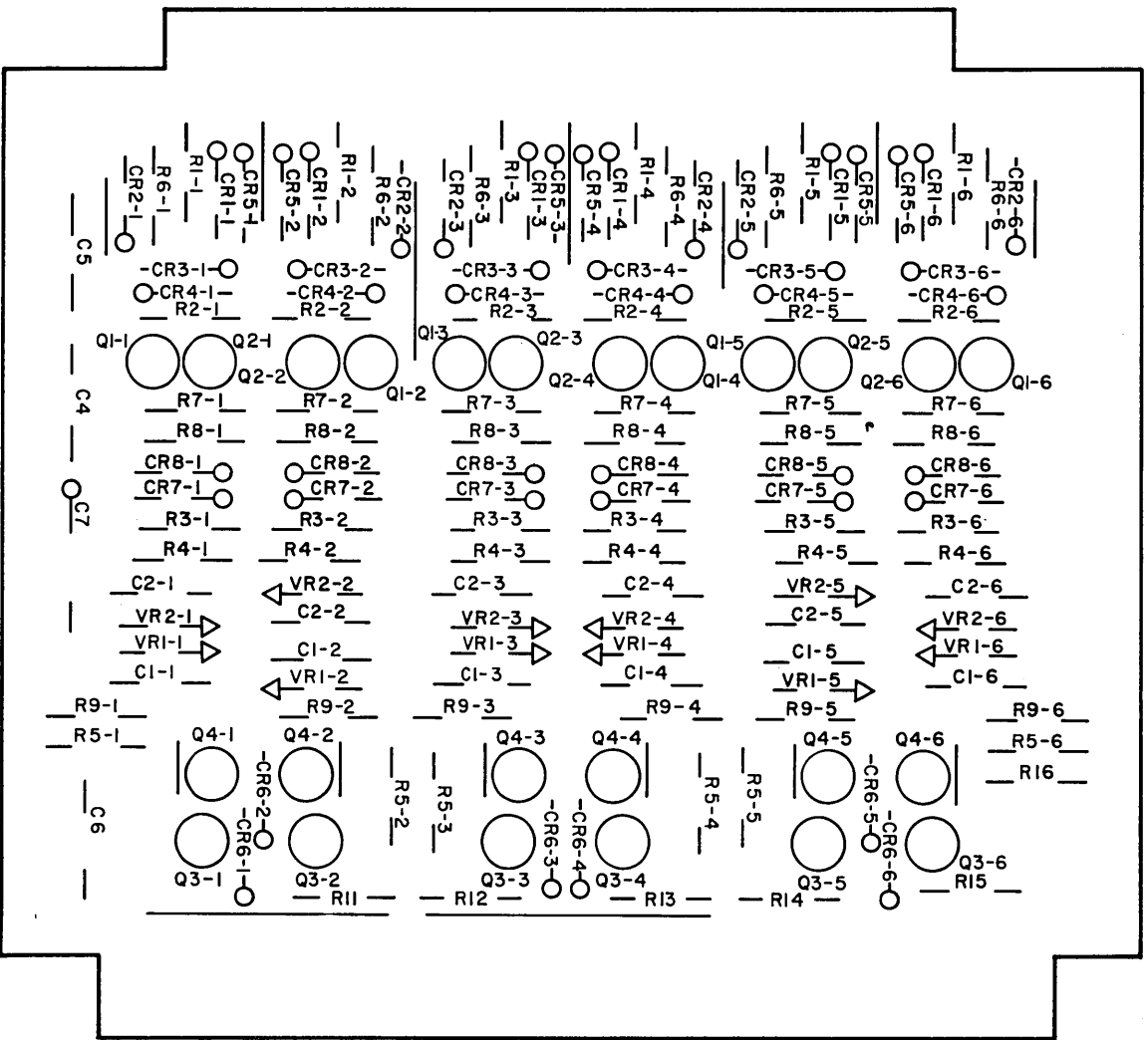
POLARIZING PINS: 32 & 46



CIRCUITS

1	43	40
2	43	38
3	43	24
4	22	23
5	22	8
6	22	7

42	41
34	35
28	26
20	11
12	10
4	5



Pin 1

Connector
End

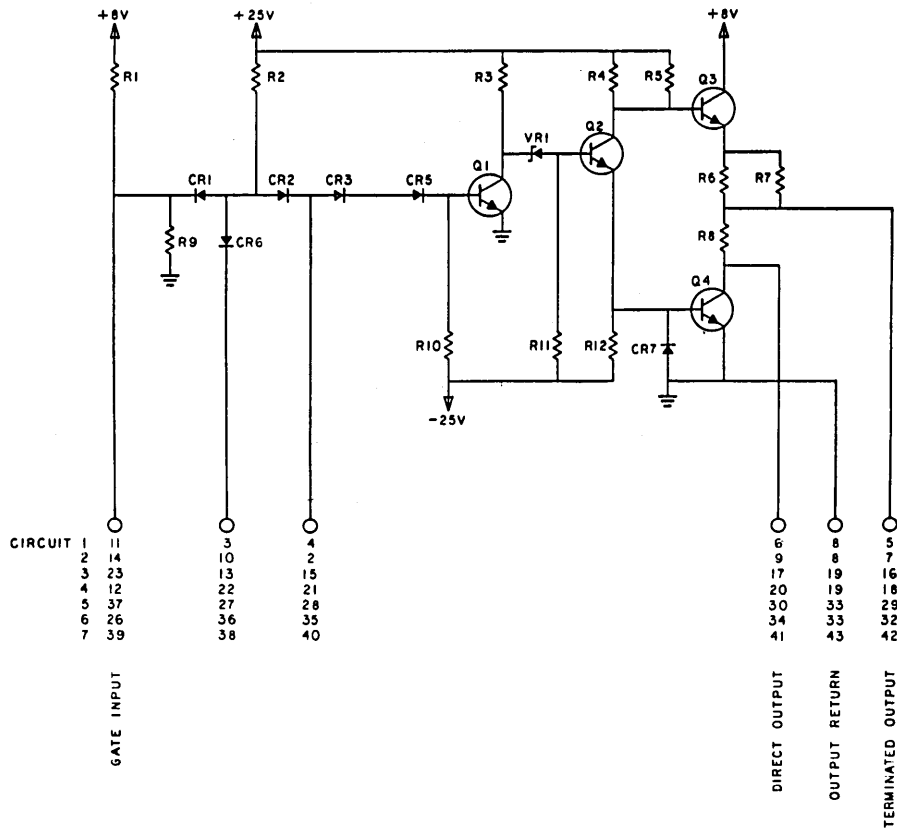
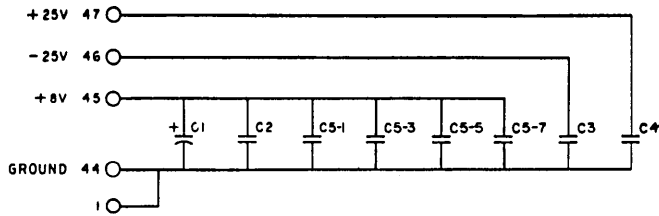
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Cable Driver No. 2 AB55
104840E

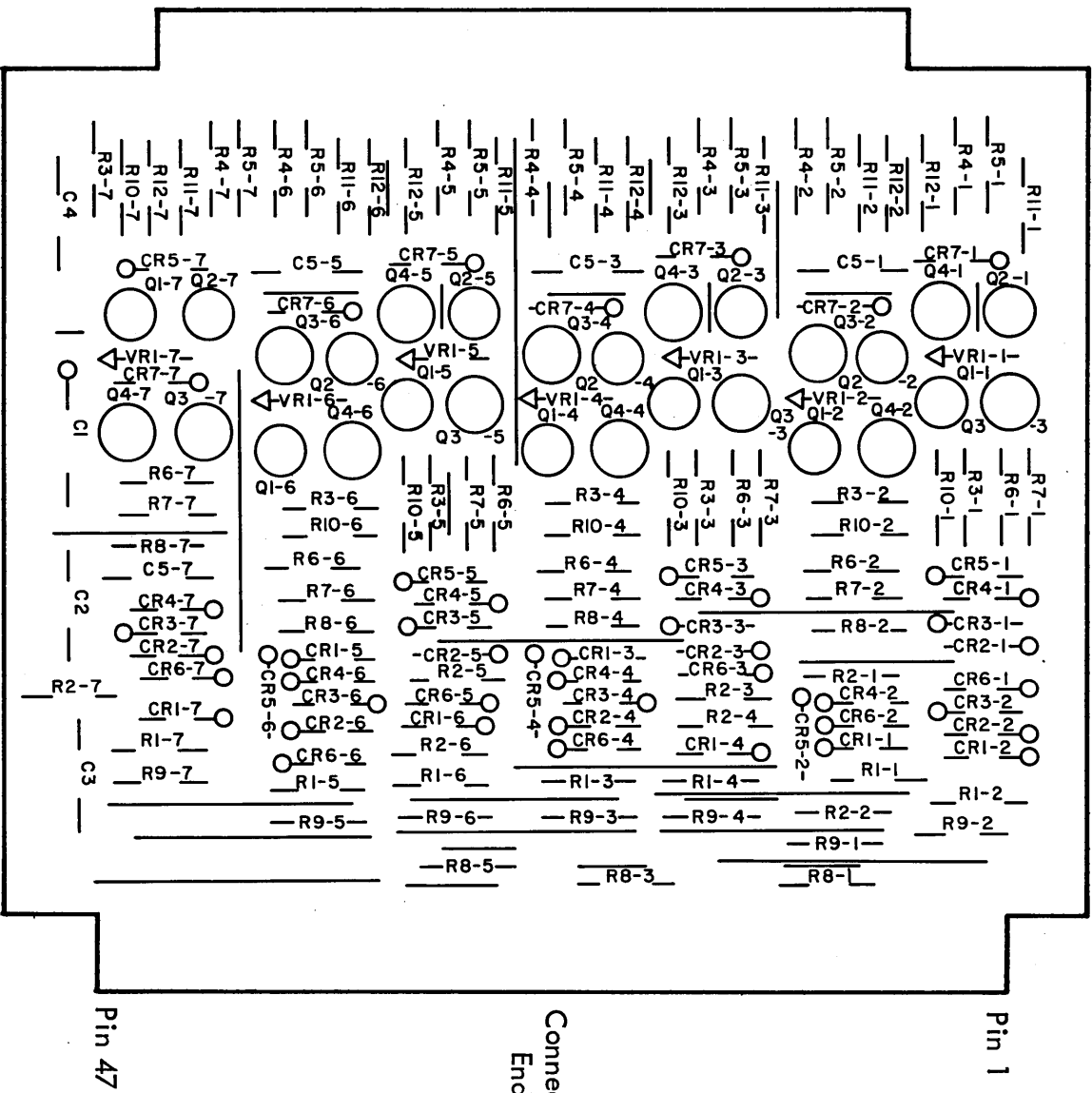
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B	SEE REV. E.O.

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EXTENSION	106315	
SCALE	D	
DATE	106315	
BY	B	

POLARIZING PINS 34 & 40

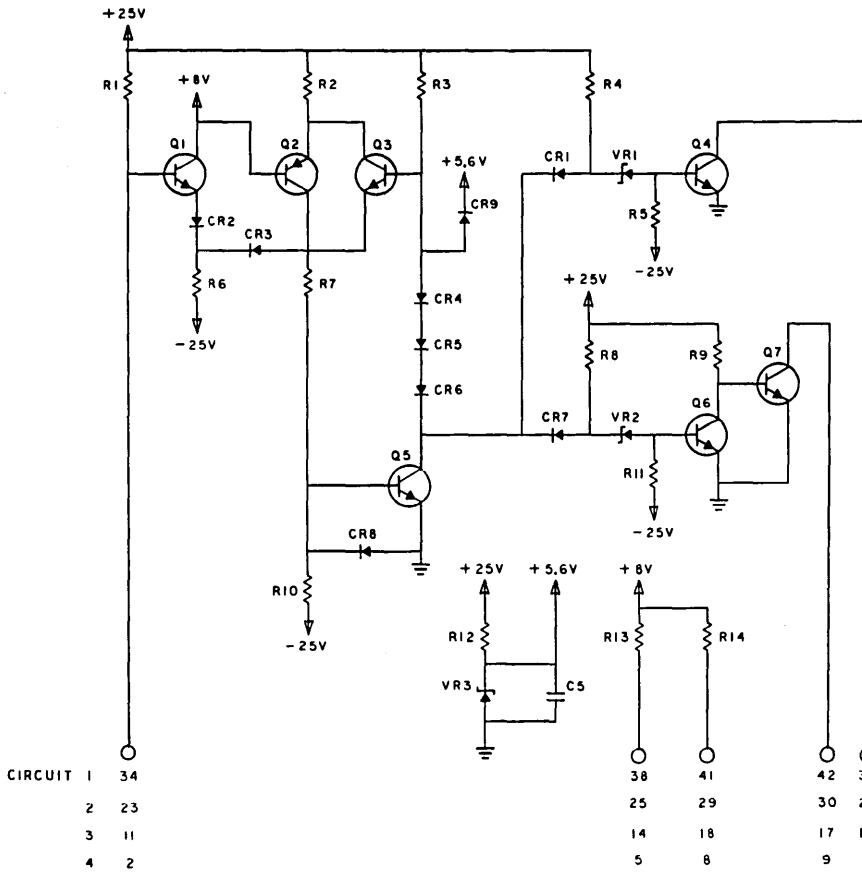
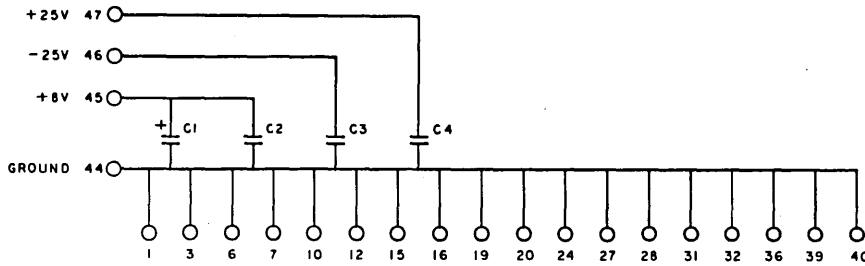


Cable Driver AK53
106317D



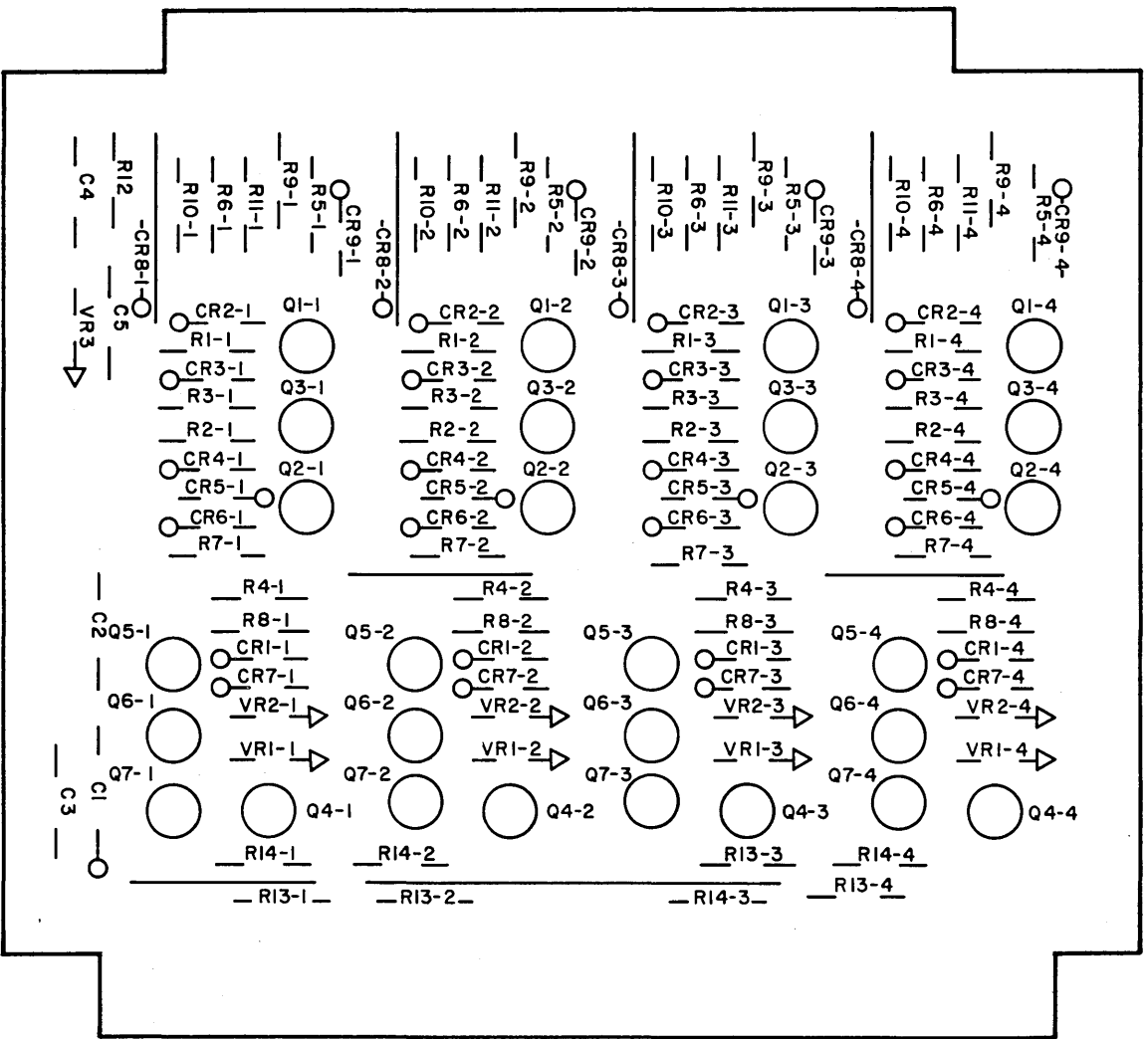
LAJ - RELEASED TO NARS

POLARIZING PINS: 3 & 31



CIRCUIT 1	34	38	41	42	37
2	23	25	29	30	26
3	11	14	18	17	13
4	2	5	8	9	4

REVISIONS		DESCRIPTION		REFERENCE DESIGNATION	
NO.	DATE	BY	CHKD.	FIG. NO.	REV.
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APPROVED FOR RELEASE AUTHORITY: 105-10449 DATE: 11-15-2013 BY: [Signature] FOR OFFICIAL USE ONLY					
SCHEMATIC SCHEMATIC NUMBER 107232					
DRAWING NO. AK54 PART NO. 107234					



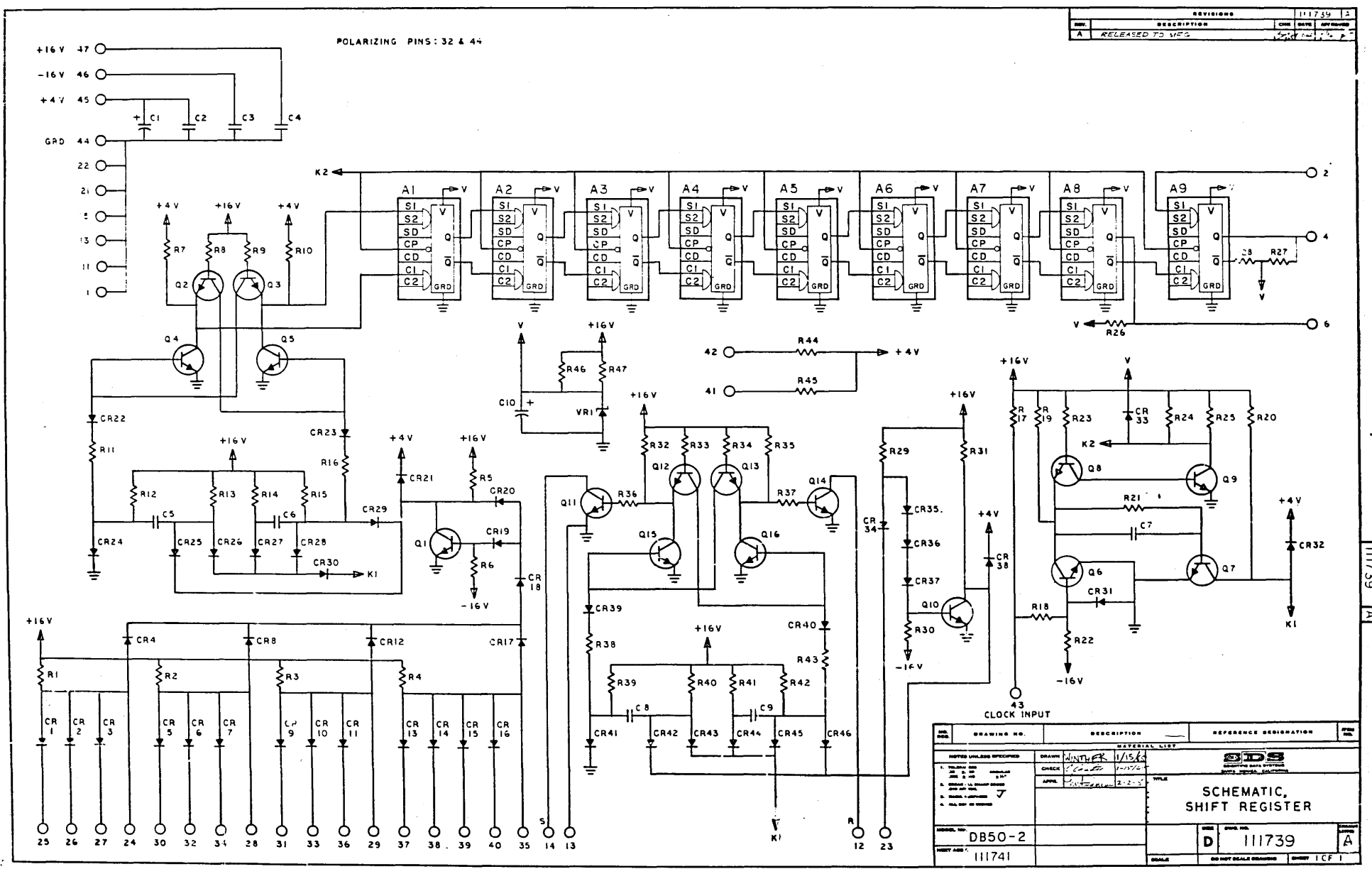
Pin 1

Connector
End

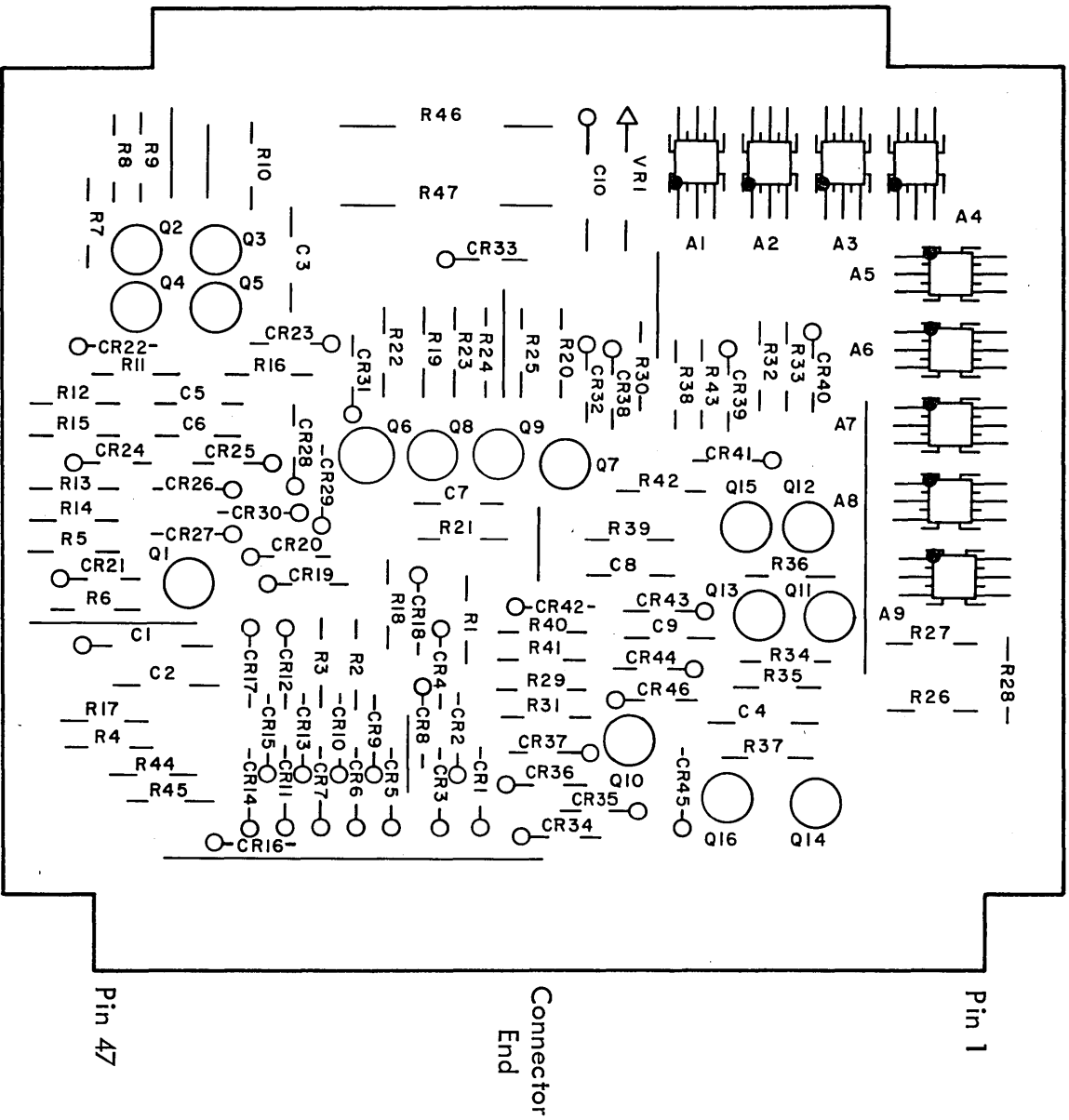
Pin 47

Schmitt Trigger AK54
107234A

REVISIONS		111739	2
REV.	DESCRIPTION	CHKD	DATE
A	RELEASED TO VFC		11-1-64



Shift Register DB50-2
111741B



Pin 1

Connector
End

Pin 47

REV. B
 DRAWING NO. 111741
 ML



MATERIAL LIST

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	111741	E

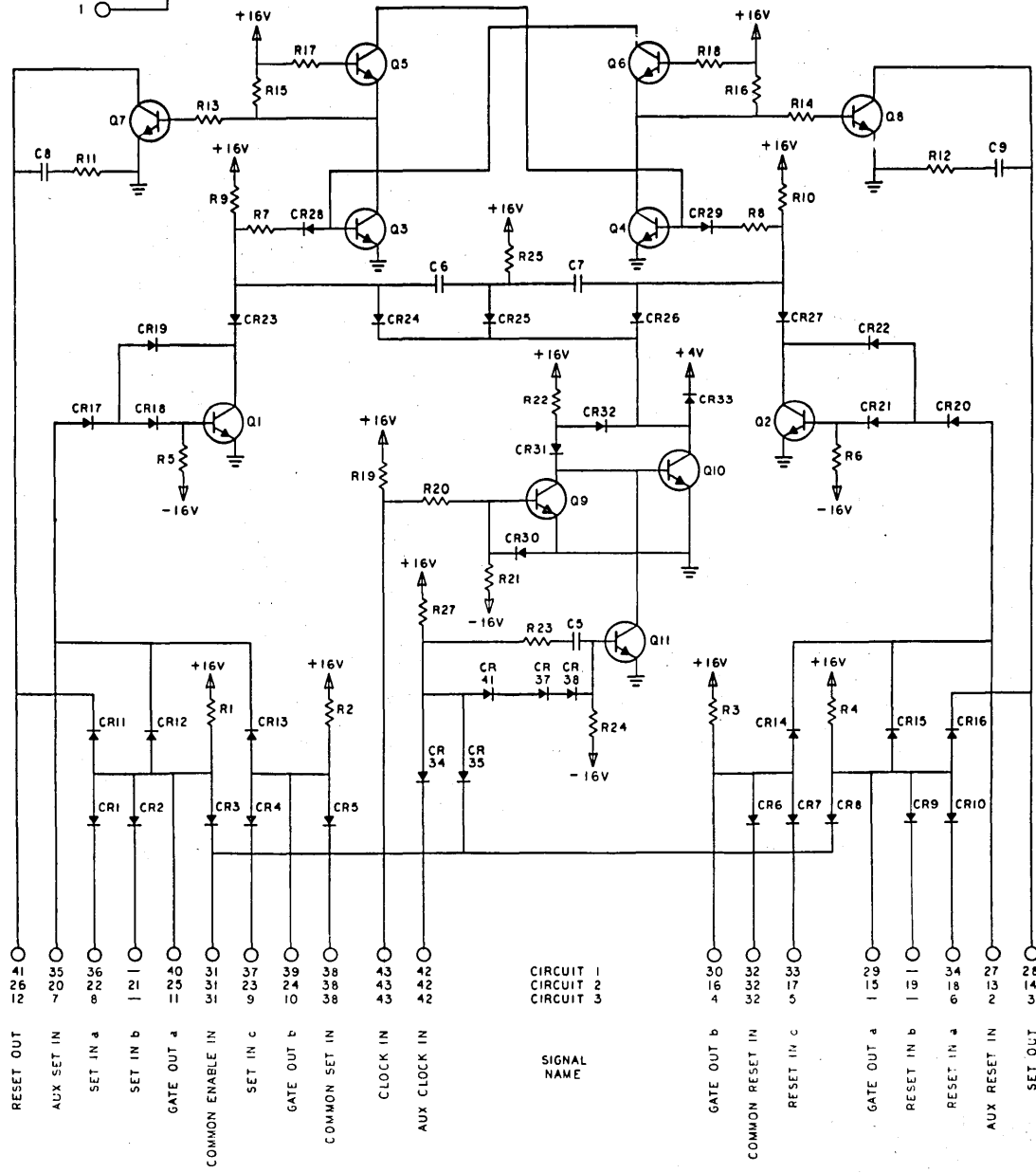
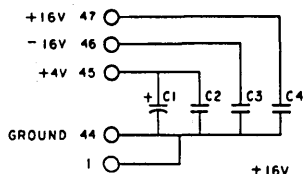
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1	Board, Printed Wiring	111740	1	
2	Handle, Circuit Card	100016	1	
3	Eyelet, Tubular	103896-016	2	
4	Strip, Marker	100197	1	
5	Contact, Conn. Upper	100097	23	
6	Contact, Conn. Lower	100098	24	
7	Transistor, SDS 216	103242	10	Q1, 4, 5, 6, 9, 10, 11, 14, 15, 16
8	Transistor, SDS 220	106781	5	Q2, 3, 12, 13, 8
9	Diode, SDS 103	100091	46	CR1 thru CR46
10	Integrated Ckt, SDS 301	108217	9	A1 thru A9
11	Resistor, 1/2 Watt	100111-332	5	R1, 2, 3, 4, 29
12	Resistor, 1/2 Watt	100111-222	5	R5, 26, 31, 27, 28
13	Resistor, 1/2 Watt	100111-153	2	R6, 30
14	Resistor, 1/2 Watt	100111-102	8	R7, 10, 13, 14, 24, 25, 40, 41
15	Resistor, 1/2 Watt	100111-562	4	R8, 9, 33, 34
16	Resistor, 1/2 Watt	100111-101	5	R11, 16, 38, 43, 21
17	Resistor, 1/2 Watt	100111-822	4	R12, 15, 39, 42
18	Resistor, 1/2 Watt	100111-122	2	R20, 19
19	Resistor, 1/2 Watt	100111-182	1	R17
20	Resistor, 1/2 Watt	100111-681	1	R18
21	Resistor, 1/2 Watt	100111-103	1	R22
22				
23	Resistor, 1/2 Watt	100111-302	3	R32, 35, 23
24	Resistor, 1/2 Watt	100111-470	2	R36, 37
25	Resistor, 1/2 Watt	100111-151	2	R44, 45
26	Capacitor, Tantalum	100312-156	2	C1, 10
27	Capacitor, Mylar	100308-103	3	C2, 3, 4
28	Capacitor, Mica	100107-820	4	C5, 6, 8, 9
29	Capacitor, Mica	100107-221	1	C7
30				
31				
32				

RELEASED TO MFG
SEE REV. 1
M.M. 10/12/62

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DATE	SCALE	ISSUE
10/12/62		D
DESIGNER	CHECKED	APPROVED
WORKS UNDER SUPERVISION	DATE	SCALE
PROJECT NO.	PROJECT NAME	SCALE
FB52	104203	

POLARIZING PINS: 28 & 34

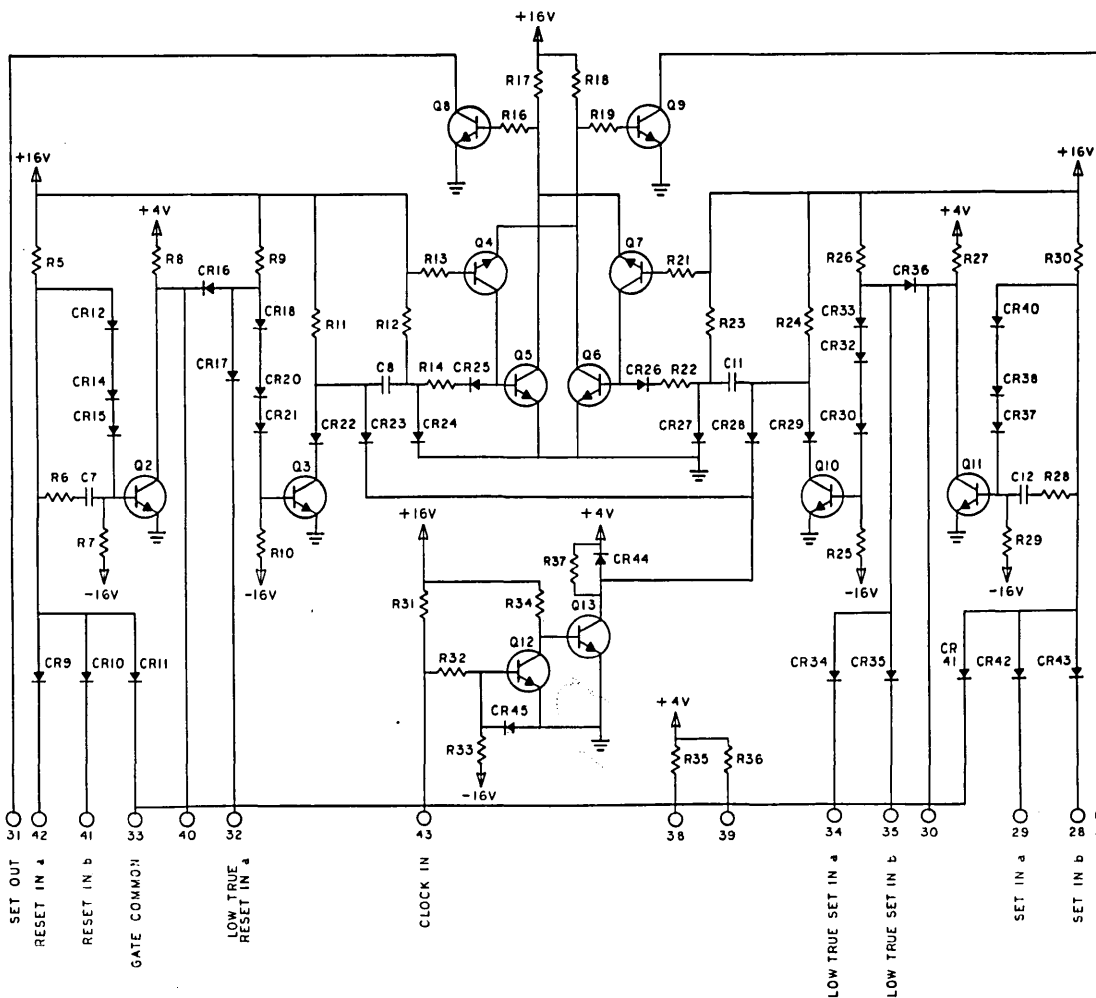
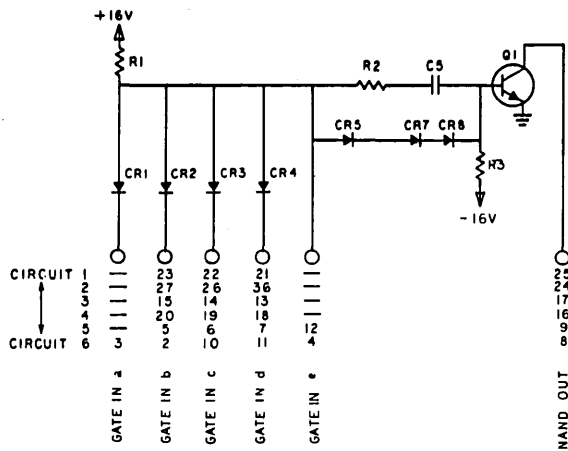
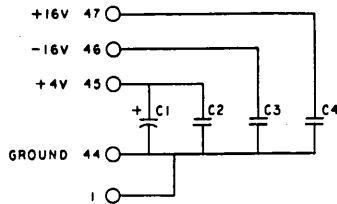


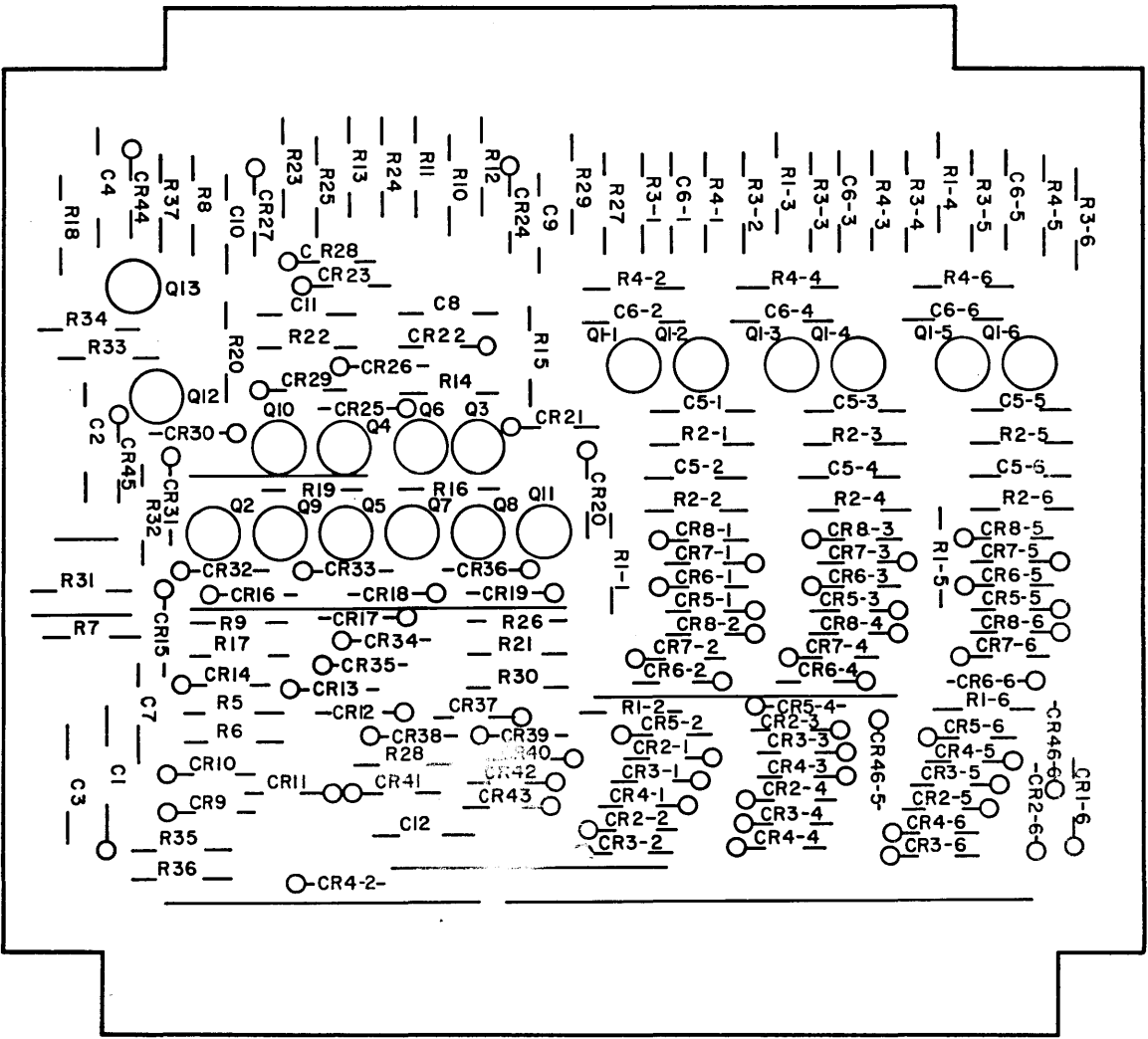
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- AUX SET IN 35
- SET IN a 36
- SET IN b 21
- GATE OUT a 40
- COMMON ENABLE IN 31
- SET IN c 37
- GATE OUT b 39
- COMMON SET IN 38
- CLOCK IN 43
- AUX CLOCK IN 42
- GATE OUT b 30
- COMMON RESET IN 16
- RESET IN c 32
- GATE OUT a 15
- RESET IN b 19
- RESET IN a 34
- AUX RESET IN 27
- SET OUT 28

REV	DESCRIPTION	DATE
A	RELEASED TO MFG.	1/14/54
B	SEE REV E.O.	1/14/54
C	SEE REV E.O.	1/14/54

DESIGNER	DESCRIPTION	DATE
W. J. ...	SCHEMATIC NAND FLIP FLOP	1/14/54
CHECKED	BY	
W. J. ...	W. J. ...	
DATE	SCALE	
1/14/54		
PROJECT NO.	104561	
WORK CENTER	FB54	
104561	104561	

POLARIZING PINS: 32 & 38





Pin 1

Connector
End

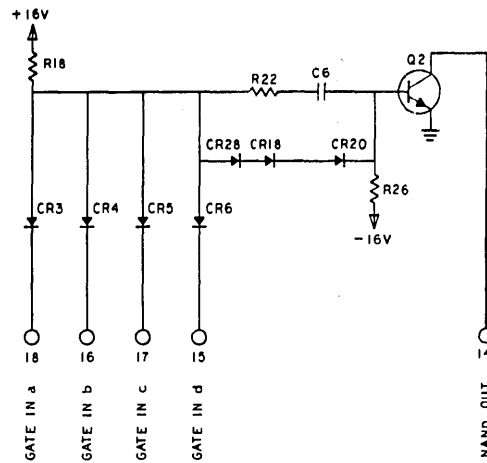
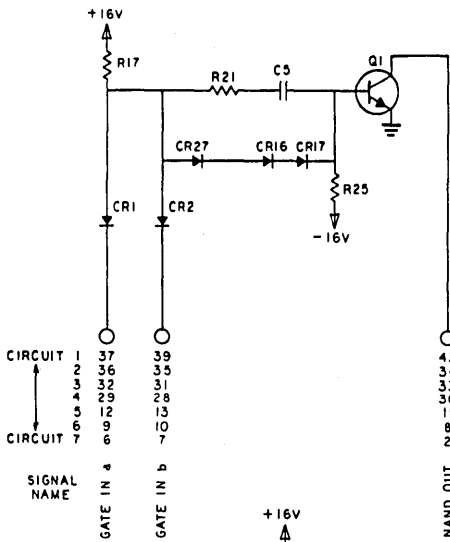
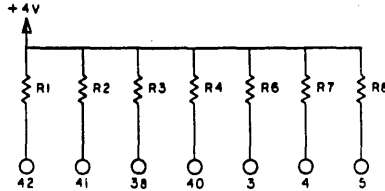
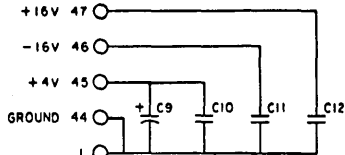
Pin 47

NAND Flip-Flop FB54
104563E

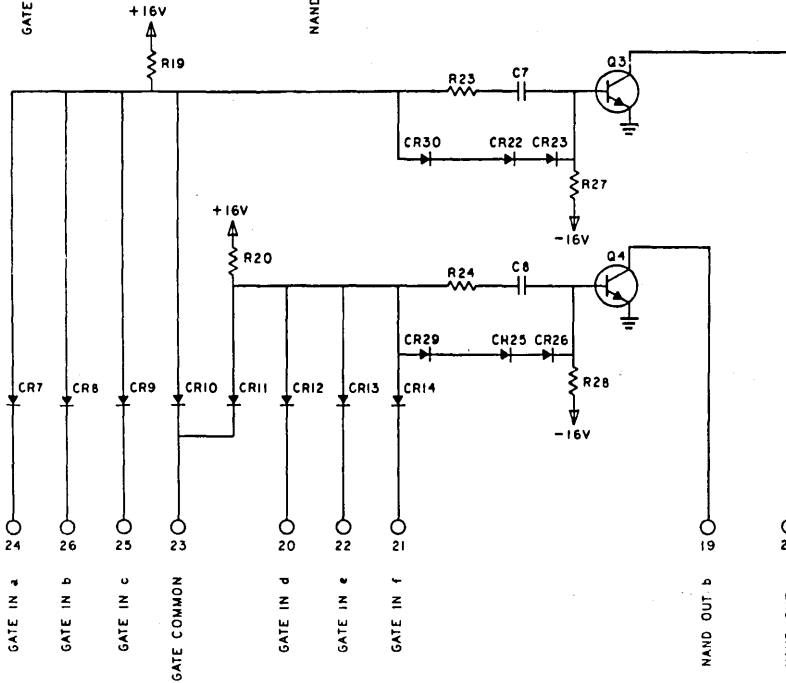
REV	DATE	BY	REVISIONS
A			REWORKED TO MFG
B			SEE REV E.O.

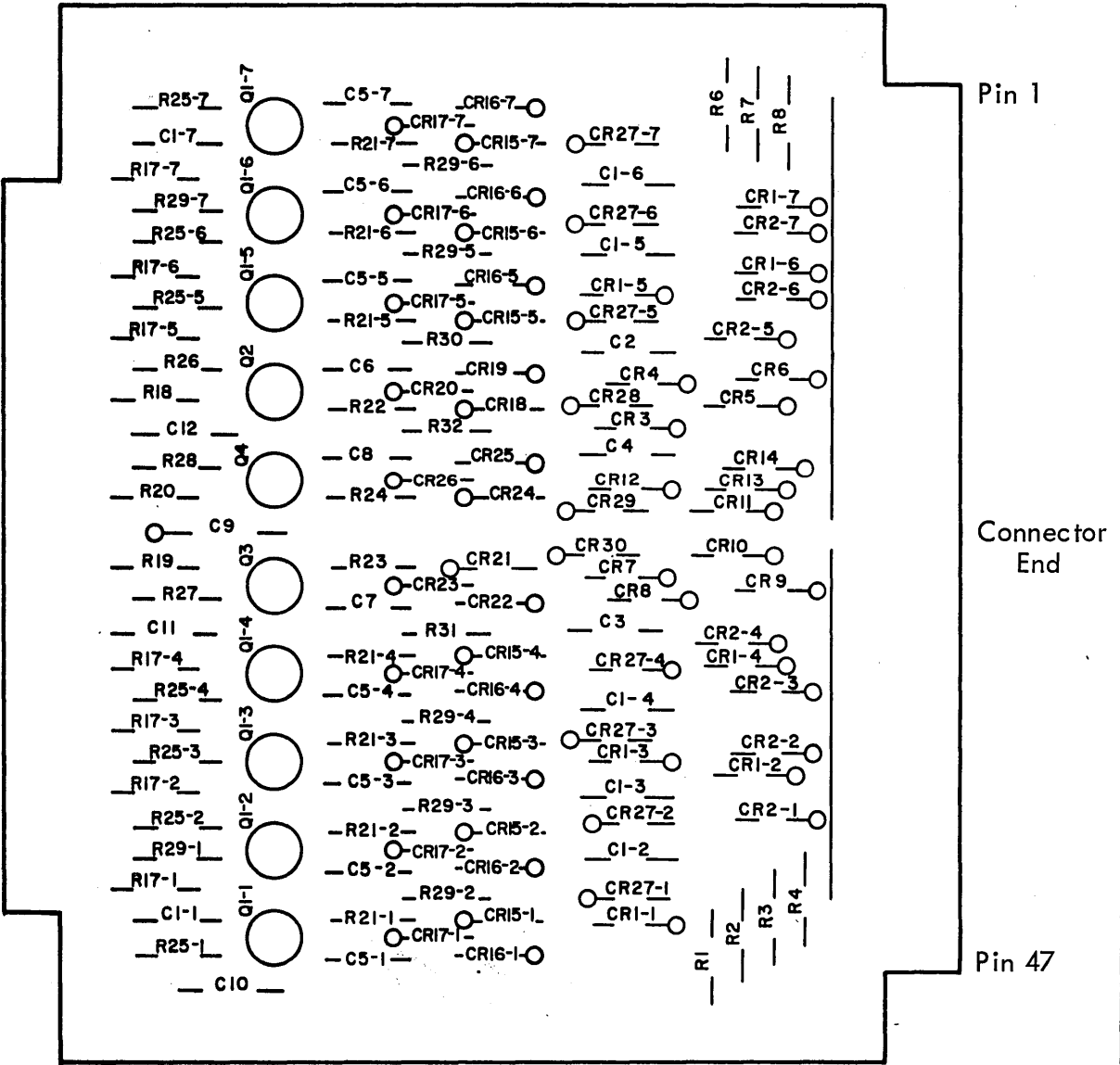
FORM NO.	DESCRIPTION	REFERENCE DESIGNATOR
103882	SCHEMATIC, NAND MODULE	
DATE	DESIGNER	CHECKER
10/24/54	W. J. WILSON	J. W. WILSON
NO. OF SHEETS	TOTAL NO. OF SHEETS	APPROVED
7	7	
DATE	NO.	NO.
10/24/54	1852	103882
REV	DATE	BY

POLARIZING PINS: 12 & 32



CIRCUIT	1	2	3	4	5	6	7
SIGNAL NAME	GATE IN a	GATE IN b	GATE IN c	GATE IN d	NAND OUT b	NAND OUT 4	

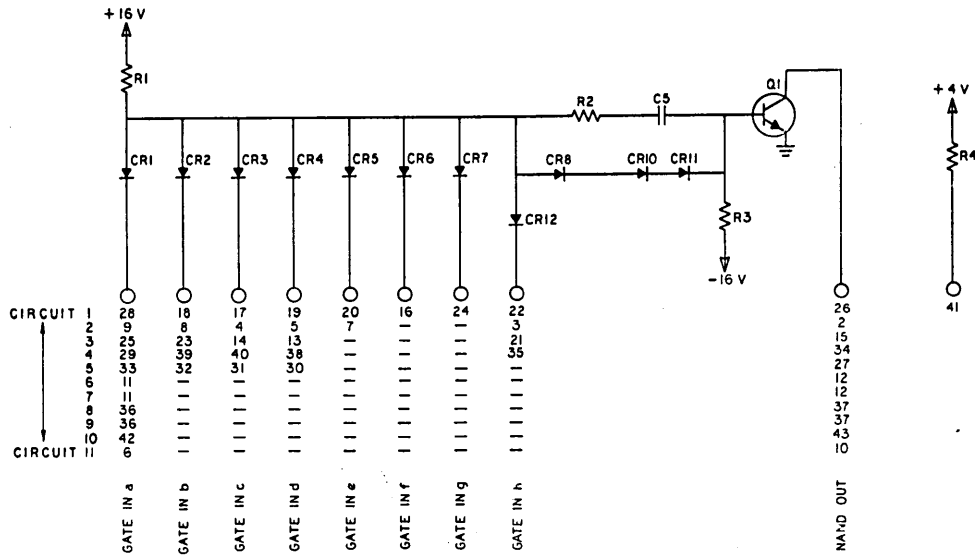
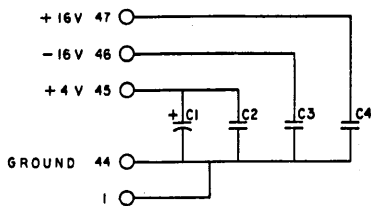




NAND Module IB52
103884D

RELEASED TO WPA
SEE REV E.O.
FORM 104554-1
REV 12-25-54

POLARIZING PINS: 32 & 36



CIRCUIT I
↑
28
9
3
4
5
6
7
8
9
10
↓
CIRCUIT II

REV	DESCRIPTION	DATE
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2	REVISION	1/15/54
3	REVISION	1/15/54
4	REVISION	1/15/54
5	REVISION	1/15/54
6	REVISION	1/15/54
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8	REVISION	1/15/54
9	REVISION	1/15/54
10	REVISION	1/15/54

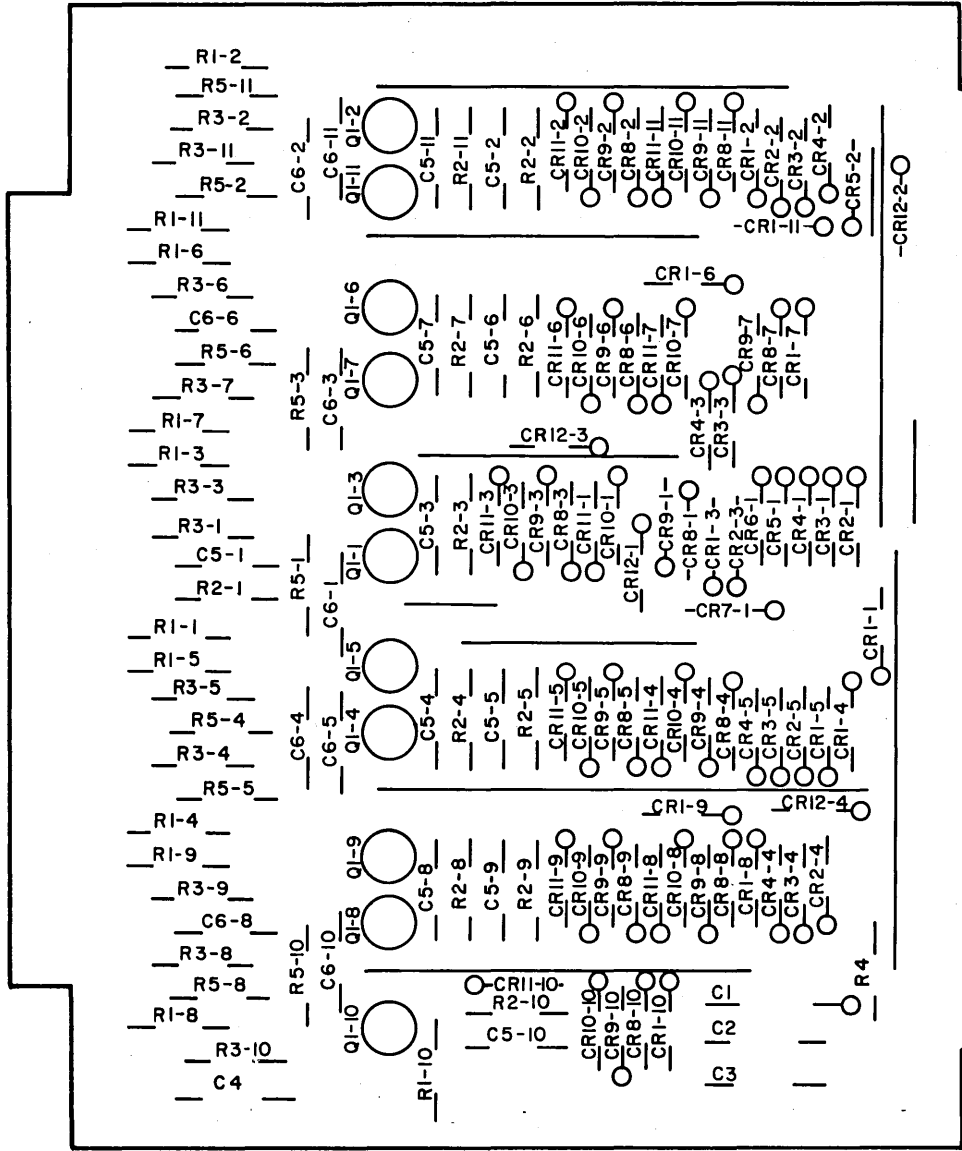
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8	REVISION	1/15/54
9	REVISION	1/15/54
10	REVISION	1/15/54

REV	DESCRIPTION	DATE
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2	REVISION	1/15/54
3	REVISION	1/15/54
4	REVISION	1/15/54
5	REVISION	1/15/54
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7	REVISION	1/15/54
8	REVISION	1/15/54
9	REVISION	1/15/54
10	REVISION	1/15/54

REV	DESCRIPTION	DATE
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3	REVISION	1/15/54
4	REVISION	1/15/54
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SCHEMATIC,
NAND #2

REV	DESCRIPTION	DATE
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6	REVISION	1/15/54
7	REVISION	1/15/54
8	REVISION	1/15/54
9	REVISION	1/15/54
10	REVISION	1/15/54



Pin 1

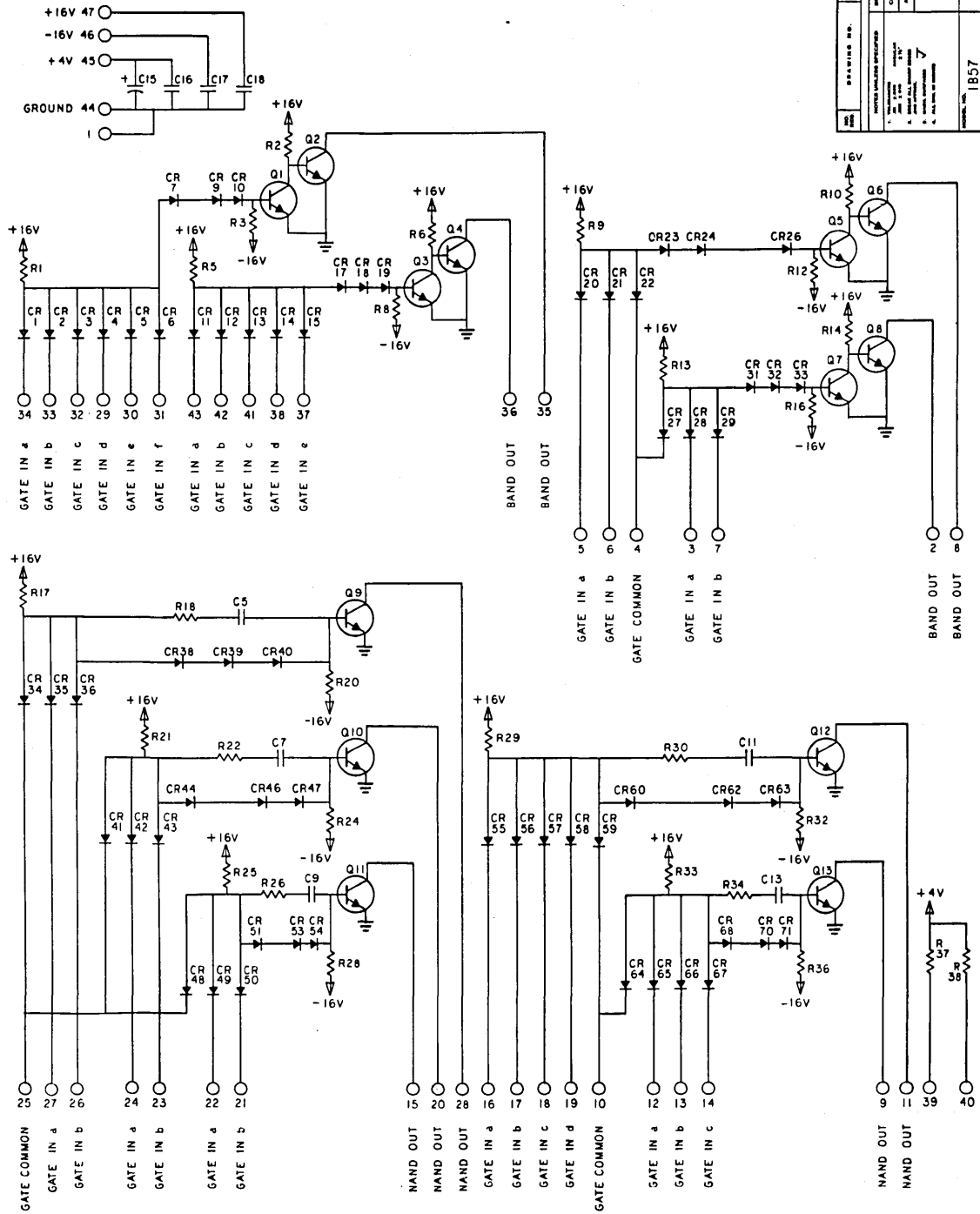
Connector End

Pin 47

NAND No. 2 IB56
104556E

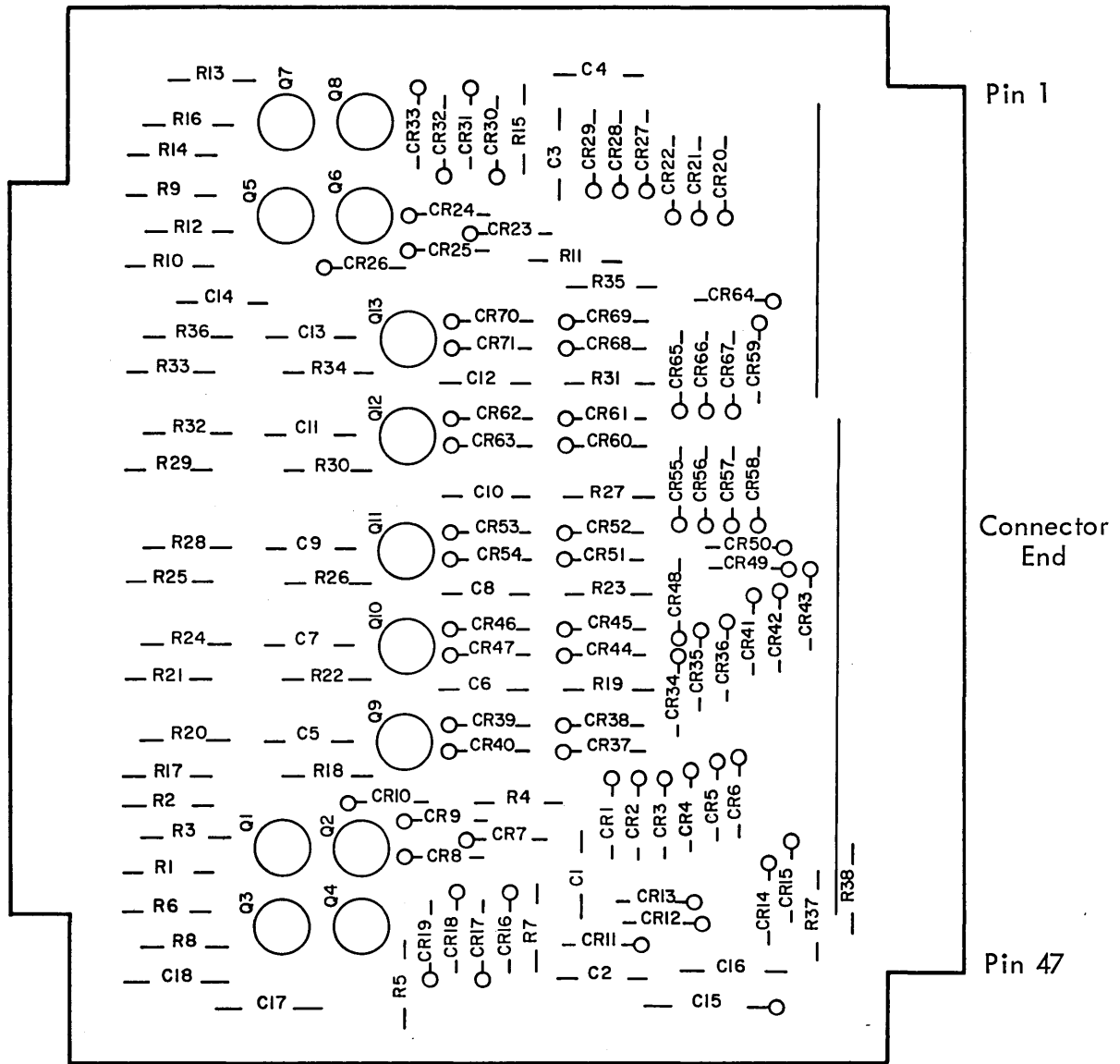
REV	DESCRIPTION	DATE
A	SEE REV E.O.	1/17
B	REVISIONS	
REVISIONS TO DRAWING		
SEE REV E.O.		

POLARIZING PINS: 30 & 44



DESIGNER	DESCRIPTION	REFERENCE NO.	REVISION
W. J. ...	SCHEMATIC LIST	104486	1 OF 1
DATE	APPROVED	SCALE	
1/17	W. J. ...		
SCHEMATIC BAND-NAND			
DRAWING NO. 104486			
REVISION 1 OF 1			

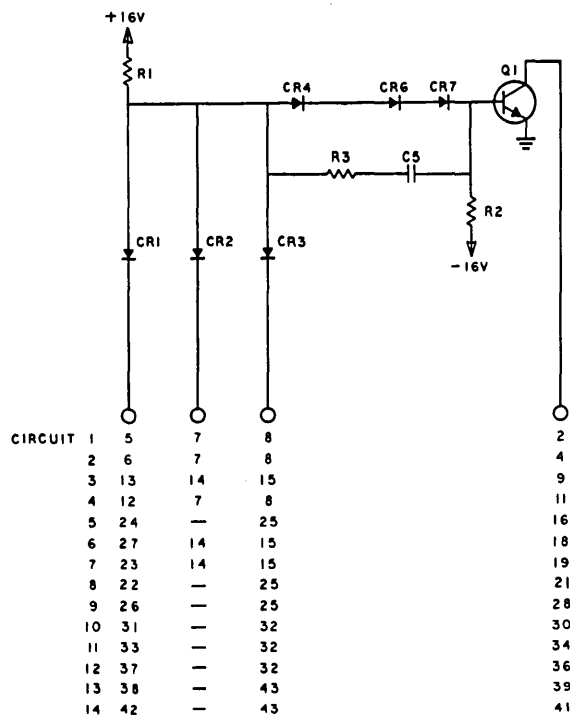
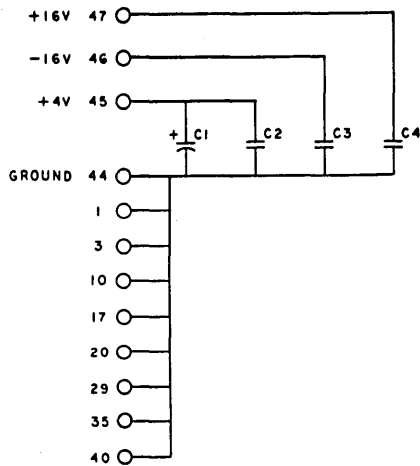
BAND NAND IB57
104488D



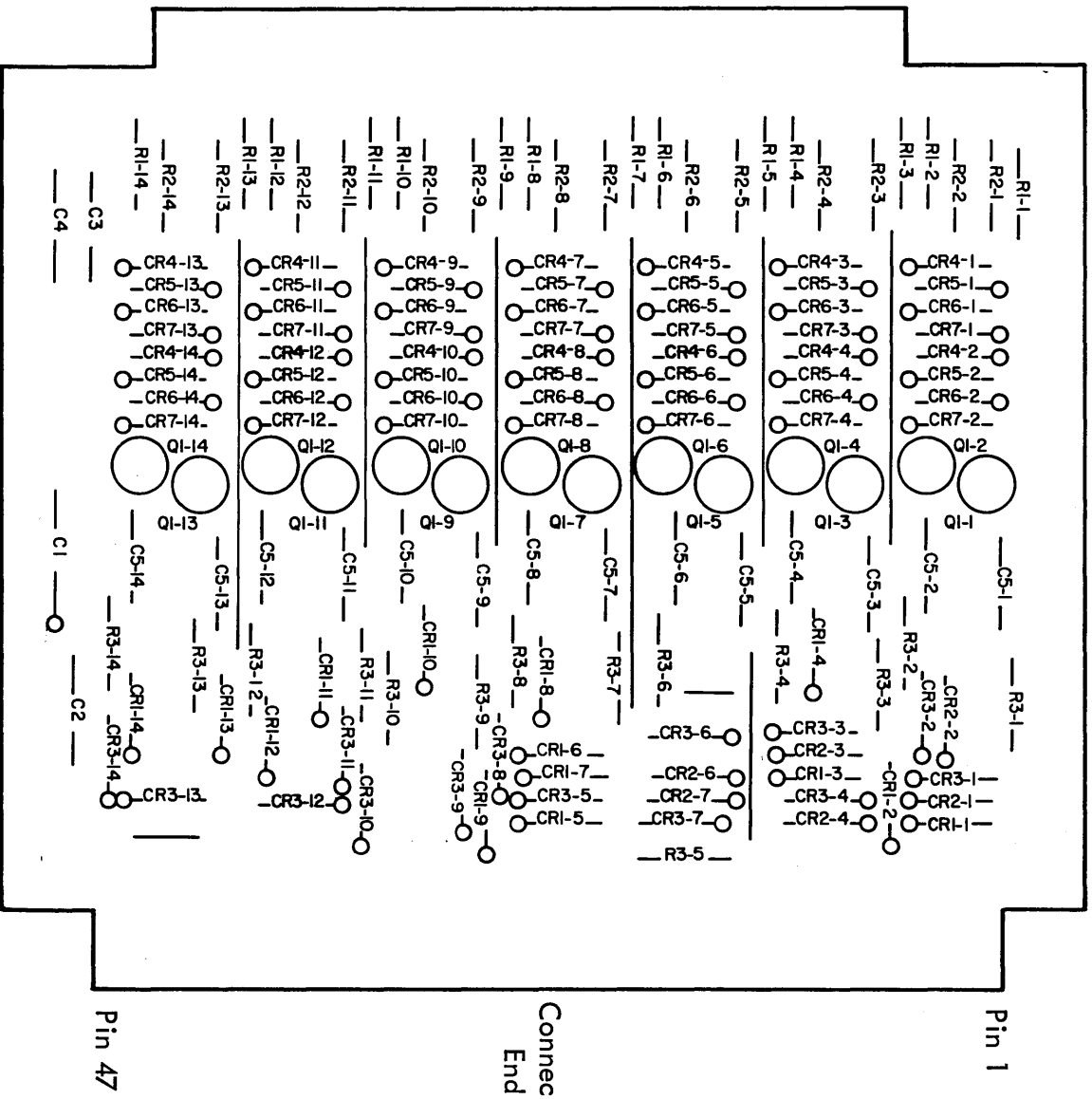
REV	D	MATERIAL LIST		SDS	SCIENTIFIC DATA SYSTEMS	ML	DWG. NO.	REV	
		DWG. TITLE					104488	0	
DWG. N		104488		MOD 0 IB57		DATE 12/4		SHEET 2 OF 2	
ML		ITEM	DWG. TITLE	DWG. NO.	NO. REQ	REMARKS ON CKT, DESIG.			
		1	Board, Printed Wiring	104487	1				
		2	Handle, Circuit Card	100016	1				
		3	Eyelet, Tubular	103896-016	2				
		4	Strip, Marker	100197	1				
		5	Contact, Conn Upper	100097	23				
		6	Contact, Conn Lower	100098	24				
		7	Transistor, SDS 216	103242	13				Q1 thru Q13
		8							
		9	Diode, SDS 103	100091	71				CR1 thru CR71
		10							
		11	Capacitor, Mica	100107-470	5				C5, 7, 9, 11, 13
		12	Capacitor, Mylar	100308-103	3				C16, 17, 18
		13	Capacitor, Tantalum	100312-156	1				C15
		14	Resistor, $\frac{1}{8}$ watt	100111-151	2				R37, 38
		15	Resistor, $\frac{1}{2}$ watt	100111-470	5				R18, 22, 26, 30, 34
		16	Resistor, $\frac{1}{2}$ watt	100111-332	9				R1, 5, 9, 13, 17, 21, 25, 29, 33,
		17	Resistor, $\frac{1}{2}$ watt	100111-153	4				R3, 8, 12, 16,
		18	Wire, Solid Bare	100042-024	10 in				
		19	Tubing, Teflon	100274-022	10 in				
		20	Resistor, $\frac{1}{2}$ watt	100111-563	5				R20, 24, 28, 32, 36
		21	Resistor, $\frac{1}{2}$ watt	100111-302	4				R2, 6, 10, 14
		22	Diode, SDS 103	100091	62				CR1 thru 7, 9 thru 15
									CR17 thru 24, 26 thru 29
									CR31 thru 36, 38 thru 44
									CR46 thru 51, 53 thru 60
									CR 62 thru 68, 70, 71

REV.	DESCRIPTION
A	REMOVED TO MEG. #1
B	SEE REV. #1

POLARIZING PINS: 42 & 46



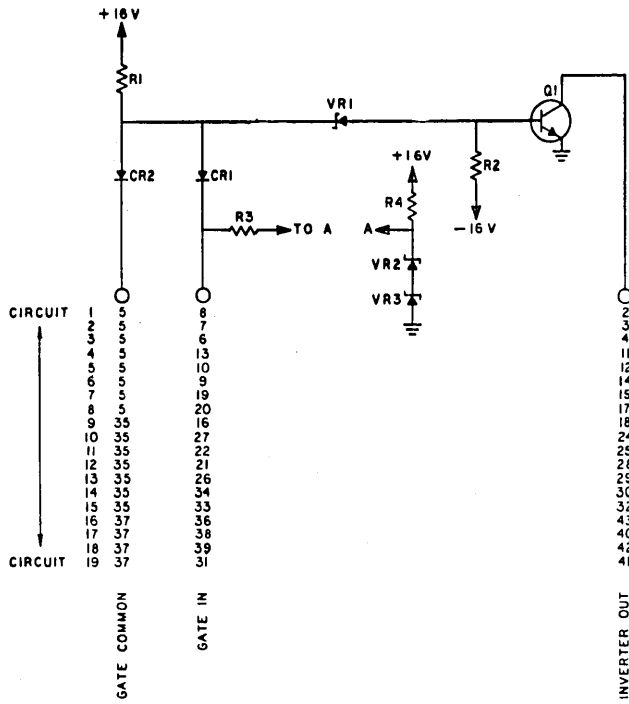
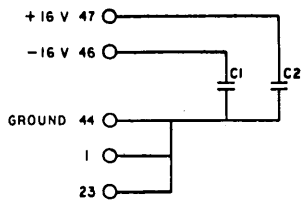
DRAWING NO.	DESCRIPTION	REFERENCE DESIGNATION
104894	SCHEMATIC, HAND # 4	B
REV.	DATE	BY
A		
B		



NAND No. 4 IB59
104896D

REV	DESCRIPTION	DATE	BY
A	RELEASED TO MFG	12/14/54	W. J. ...
B	SEE REV. E.O.	12/14/54	W. J. ...

POLARIZING PINS : 28 & 40



REV	DESCRIPTION	DATE	BY
A	RELEASED TO MFG	12/14/54	W. J. ...
B	SEE REV. E.O.	12/14/54	W. J. ...

REV	DESCRIPTION	DATE	BY
A	RELEASED TO MFG	12/14/54	W. J. ...
B	SEE REV. E.O.	12/14/54	W. J. ...

REV	DESCRIPTION	DATE	BY
A	RELEASED TO MFG	12/14/54	W. J. ...
B	SEE REV. E.O.	12/14/54	W. J. ...

REV	DESCRIPTION	DATE	BY
A	RELEASED TO MFG	12/14/54	W. J. ...
B	SEE REV. E.O.	12/14/54	W. J. ...

REV	DESCRIPTION	DATE	BY
A	RELEASED TO MFG	12/14/54	W. J. ...
B	SEE REV. E.O.	12/14/54	W. J. ...

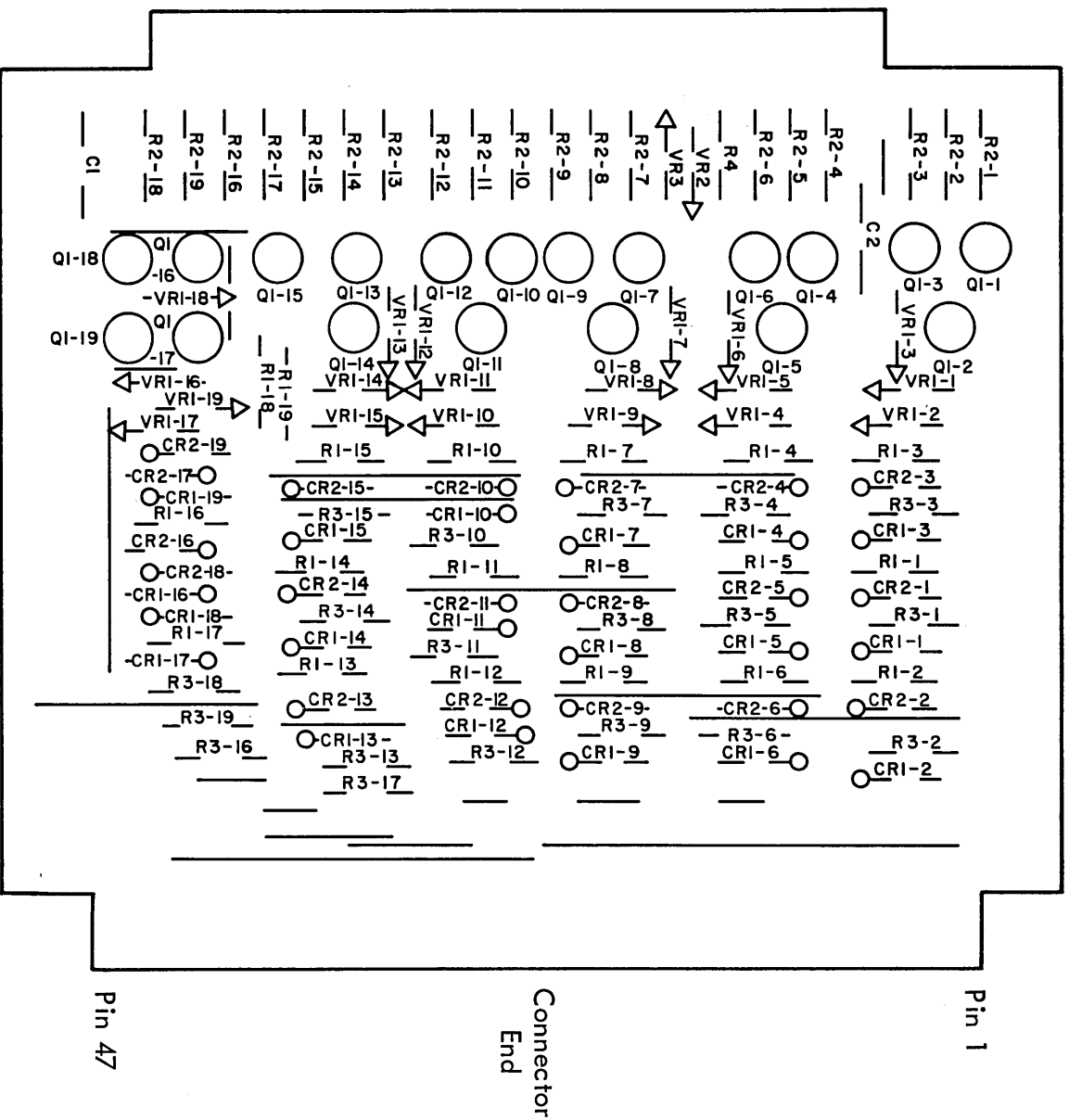


SCHEMATIC,
+ 8 TO + 4 INTERFACE

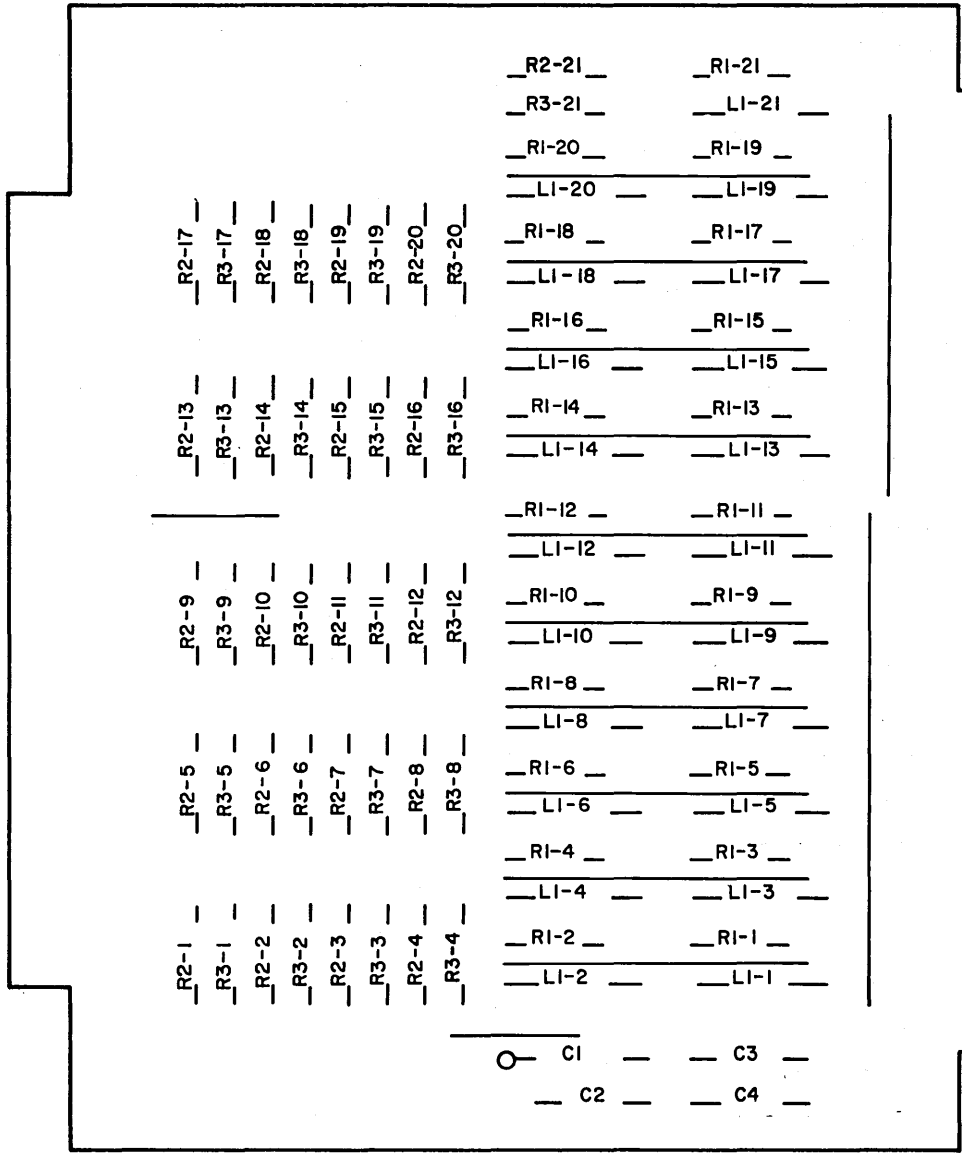
D 104641

NB50

104642



+8 To +4 Interface NB50
104642C



Pin 1

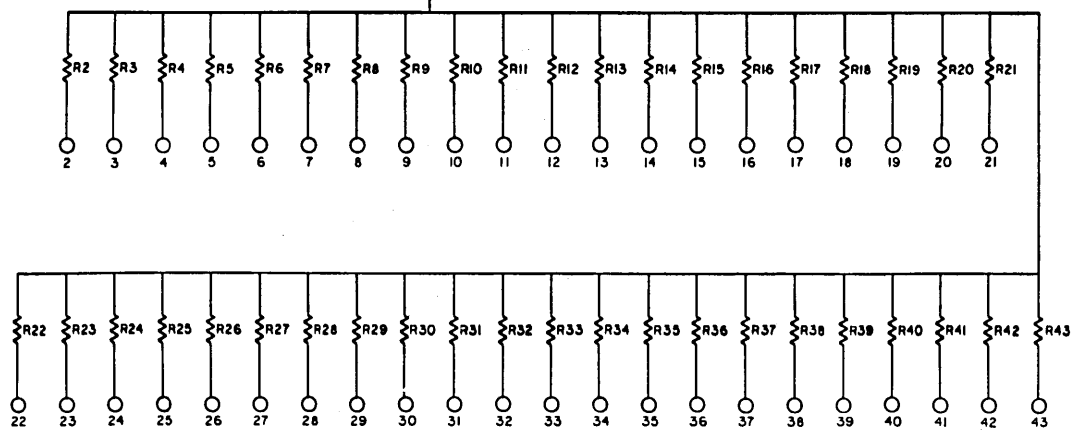
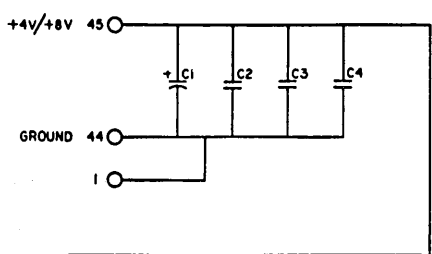
Connector End

Pin 47

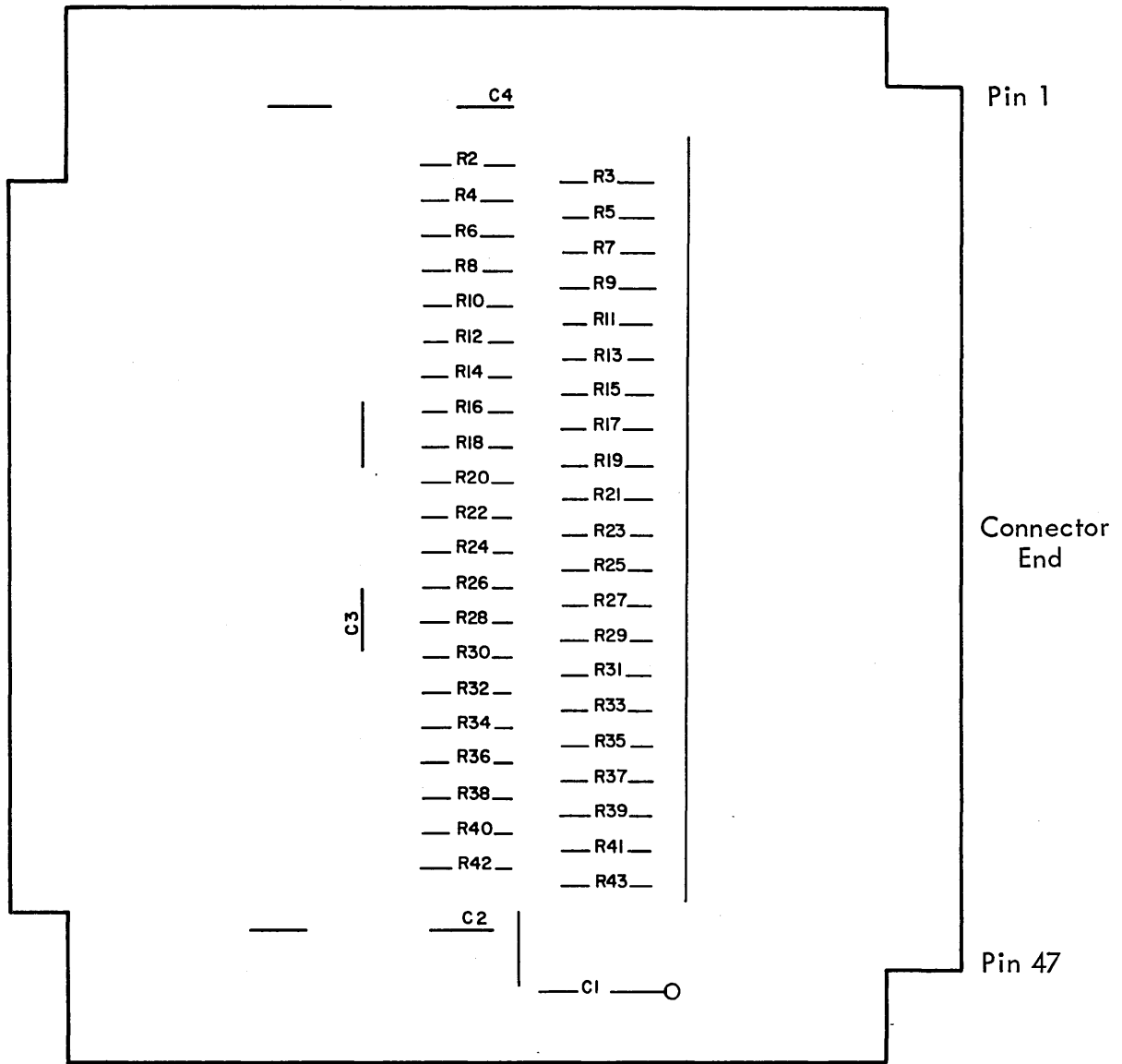
Drive Cable Interface NB52
104866B

REV	DESCRIPTION	DATE
A	RELEASED TO WPG	10/11/77

POLARIZING PINS: 12 & 40



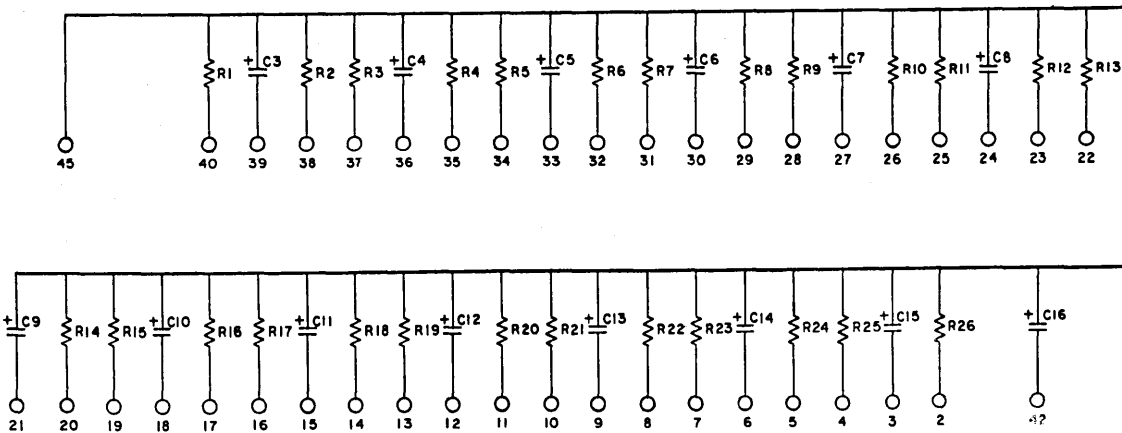
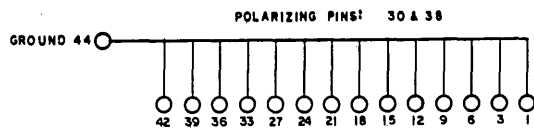
DRWG NO.	DESCRIPTION	DATE
ZB50/ZB51	SCHEMATIC, TERMINATION MODULE	10/40/76
REV	DESCRIPTION	DATE
D	104014	10/40/76

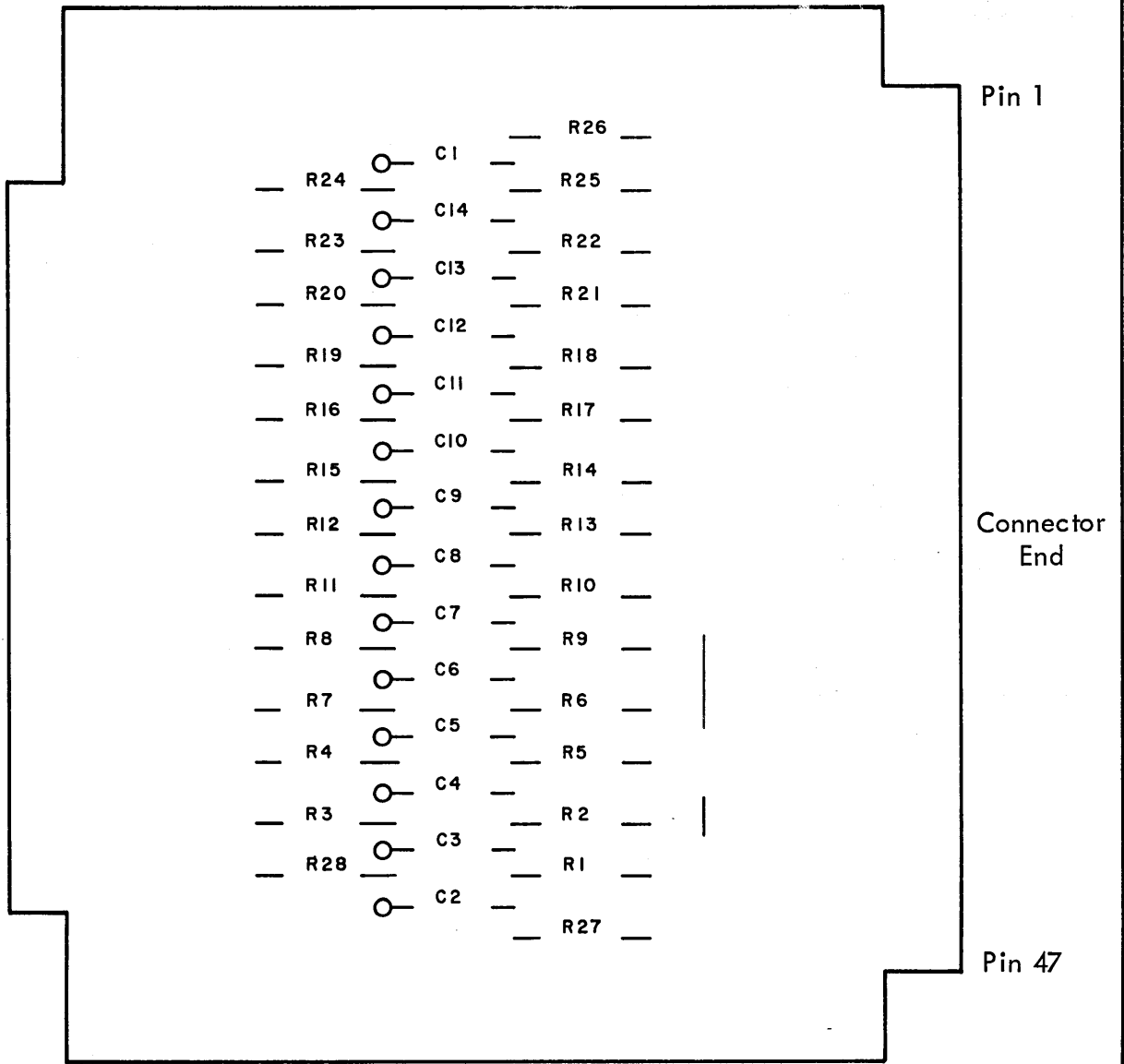


Termination Module +4 ZB50
104016A

104329 C	
REV	DESCRIPTION
A	PRE RELEASED
B	ADDED 0271528 CIG & REMOVED JUMPER FROM PIN 43 TO PIN 45
C	RELEASED TO MFG.
C. SEE REV E.O.	

DESIGN NO.	ZB52	REVISION	D
PROJECT NO.	104331	DATE	10/9/56
DESCRIPTION		SCHEMATIC, CABLE TERMINATION MODULE	
STANDARD LIST		W/W/W 10/9/56 C/MS/STP 10/15/56 C/C/L/S 1/27/56	
REFERENCE DESIGNATION		104329 C	





Cable Termination Module ZB52
104331E

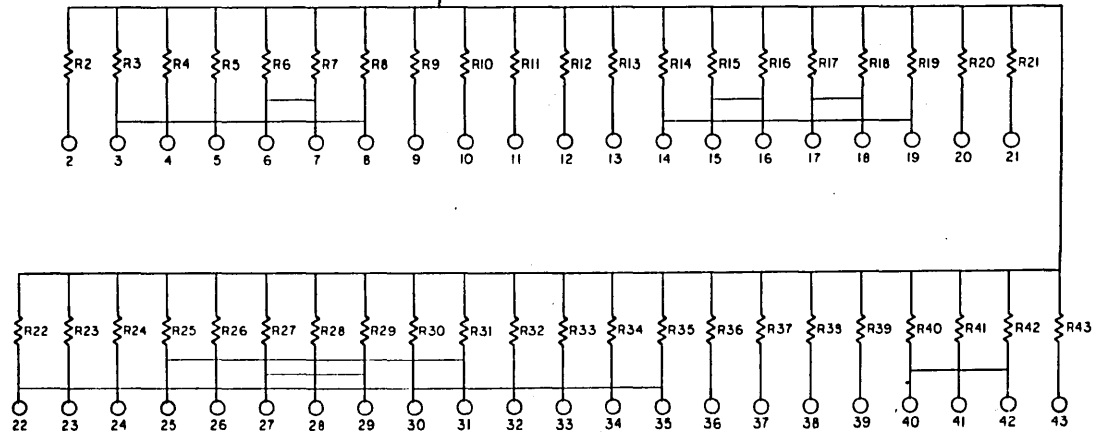
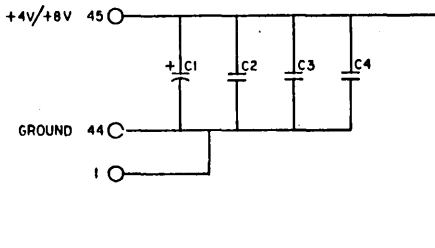
REV	DESCRIPTION	DATE	BY
A	RELEASED TO MFG		

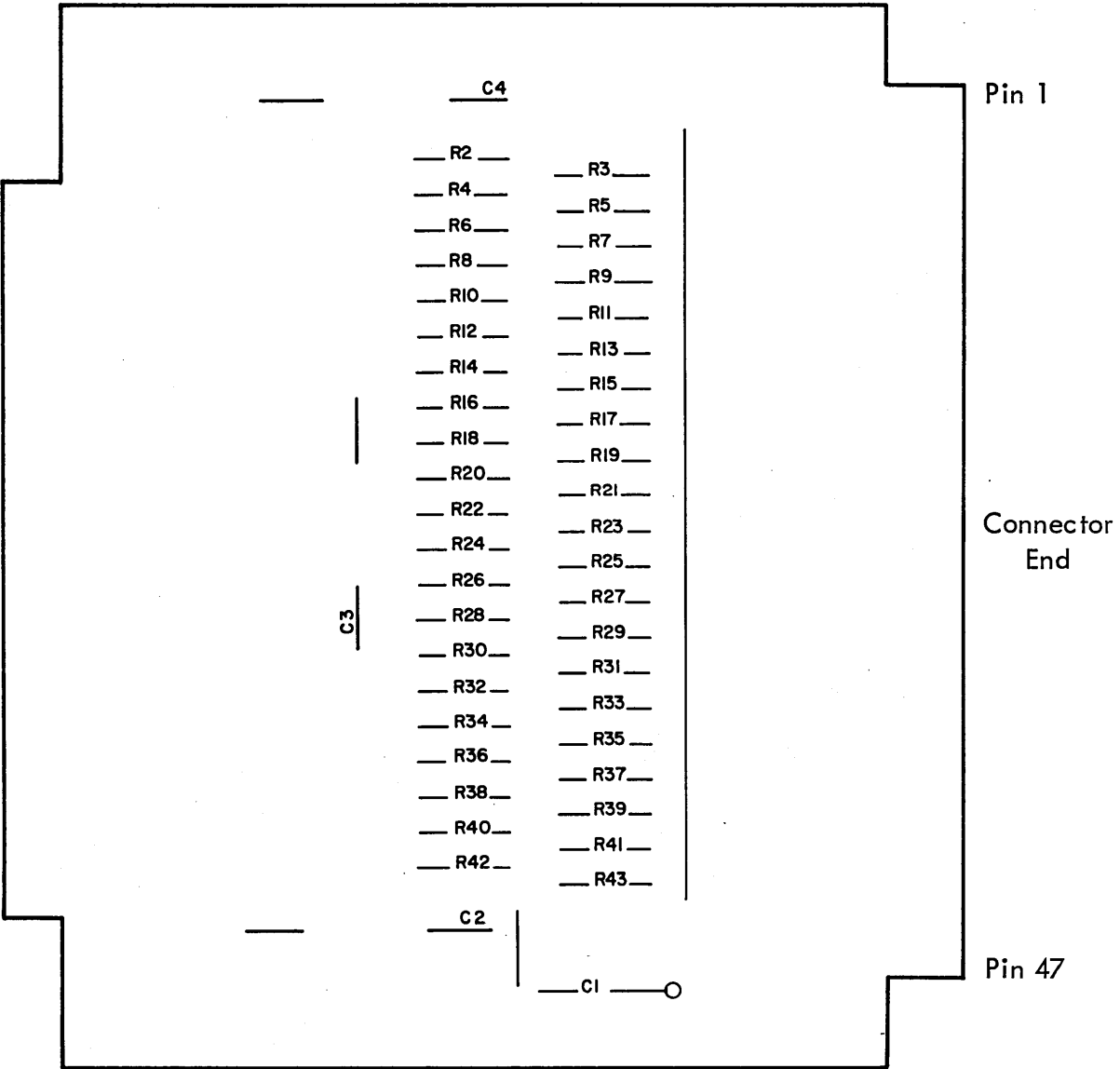
DRAWING NO.		DESCRIPTION		REFERENCE DESIGNATION	
ZB55		104993		D	
104994		SCHEMATIC, TERMINATION MODULE		A	
DATE		APPROVED		SCALE	
				1:1	

NOTES:

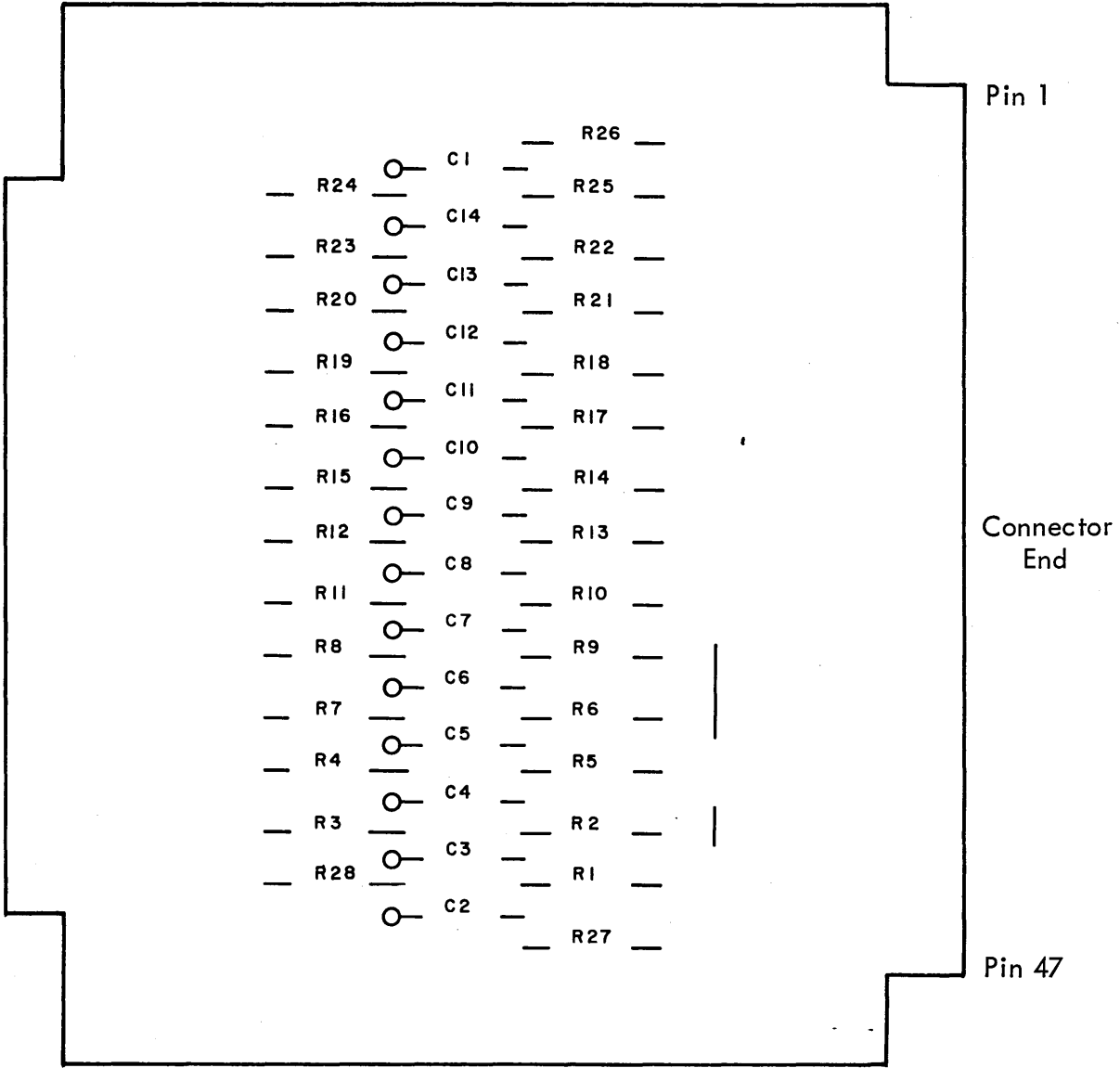
- 1. R3, R5 THRU R9, R11, R14 THRU R33, R35, R40, R41 & R42 ARE OMITTED.

POLARIZING PINS: 12 & 40






Termination Module ZB55
104994B



Cable Termination Module ZB73
111991A

REV.	A	 MATERIAL LIST	ML	DRAWING NO.	REV.
				111991	A
DRAWING NO.	111991	Assy. P. W. Cable DRAWING TITLE <u>Termination Module</u> MODEL NO. <u>ZB73</u> DATE <u>2/1/65</u> SHEET <u>2</u> OF <u>2</u>			
		ITEM NO.	DRAWING TITLE	DWG. NO.	NO. REQ.
ML	1	Board, Printed Wiring	104330	1	
	2	Handle, Circuit Card	100016	1	
	3	Eyelet, Tubular	103896-016	2	
	4	Strip, Marker	100197	1	
	5	Contact, Conn. Upper	100097	23	
	6	Contact, Conn. Lower	100098	24	
	7	Capacitor, Tantalum	100312-156	14	C1 thru C14
	8	Resistor, Metal Film	100680-330	28	R1 thru R28
	9	Wire, Solid Bare	100042-024	1in	
	10	Tubing, Teflon	100274-022	1in	
	11	Heatsink(Extruded)	106579	1	
	12	Screw, Flat Hd. Phillips ^{100°}	100012-307	2	
	13	Washer, Flat	100018-300	2	
	14	Washer, Lock Int. Tooth	100024-300	2	
	15	Nut, Hex Machine	100008-300	2	
	16	Schematic	113043	x	ref
	17	Dwg List	113044	x	ref
	18	Test Spec	113045	x	ref