

File No. S370-36  
Order No. SY20-0880-1

**Systems**

## **IBM Virtual Machine Facility/370: Control Program (CP) Program Logic**

**Release 1 PLC 9**

This publication describes the internal logic of the VM/370 control program. Major sections discuss:

- The function of the control program
- The control program's modules
- The control program's data areas

Diagnostic information is also included.

This publication is intended for IBM personnel responsible for program maintenance.

Prerequisites for a thorough understanding and for the effective use of this publication are:

*IBM System/360 Principles of Operation, GA22-6821*  
*IBM System/370 Principles of Operation, GA22-7000*  
*IBM System/360 Operating System: Assembler  
Language, GC28-6514*

For titles and abstracts of other associated publications, see the publication *IBM System/360 and System/370 Bibliography, GA22-6822*

# **IBM**

Second Edition (April 1973)

This edition, together with Technical Newsletter SN20-2624, dated August 15, 1973, applies to Release 1 PLC 9 (Program Level Change) of IBM Virtual Machine Facility/370 (VM/370) and to all subsequent releases until otherwise indicated in new editions or Technical Newsletters.

Changes are continually made to the specifications herein; before using this publication in conjunction with the operation of IBM systems, consult the latest IBM System/360 and System/370 Bibliography, Order No. GA22-6822, and the IBM System/370 Advanced Function Bibliography, Order No. GC20-1763, for editions that are applicable and current.

Changes and additions are indicated by a vertical line to the left of the change.

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A form for reader's comments is provided at the back of this publication. If the form has been removed, comments may be addressed to IBM Corporation, VM/370 Publications, 24 New England Executive Park, Burlington, Massachusetts, 01803. Comments become the property of IBM.

## PREFACE

This Program Logic Manual (PLM) is a detailed guide to the VM/370 control program. It supplements the program listings by providing descriptive text, diagrams, and flowcharts. It is intended for IBM personnel responsible for program maintenance and is to be used with the following PLMs for maintaining the System Control Program (SCP).

IBM Virtual Machine Facility/370: Conversational Monitor System (CMS), Program Logic, SY20-0881

IBM Virtual Machine Facility/370: Service Routines Program Logic, SY20-0882

IBM System/360 Operating System: Assembler (F), GY26-3700

IBM CALL/360-OS BASIC System Manual, GY20-0530

The logic described in this publication is about programs that are discussed in the following publication:

IBM Virtual Machine Facility/370: Introduction, GC20-1800

IBM Virtual Machine Facility/370: Planning and System Generation Guide, GC20-1801

IBM Virtual Machine Facility/370: Operator's Guide, GC20-1806

IBM Virtual Machine Facility/370: System Messages, GC20-1808

IBM Virtual Machine Facility/370: Command Language Guide for General Users, GC20-1804

IBM Virtual Machine Facility/370: System Programmer's Guide, GC20-1807

Information in this publication (if any) about the following is for planning purposes only:

- IBM System/370 Model 165 II

In this publication, the term "3330 series" is used in reference to the IBM 3330 Disk Storage, Models 1 and 2 and the IBM 3333 Disk Storage and Control, Model 1.

## MANUAL ORGANIZATION

This publication is divided into seven sections:

- The "Introduction" presents a general discussion about the functions and program organization of the VM/370 control program.
- The section "Method of Operation" contains a detailed discussion about the functions of the control program.
- "Program Organization" contains the flowcharts.
- The "Directories" contain an alphabetical list of all the CP labels that are discussed within this manual. Accompanying the label is a brief description of the function or for the subroutines, the name of the module that it is in, and cross-reference to its location in the PLM.
- The section "Data Areas" contains a detailed description of the control program data areas.
- "Diagnostic Aids" contains cross-reference information about commands, messages, wait codes, and ABEND codes.
- The Appendixes contain coding conventions, system equates, and DASD record formats.



### Addition of the Following VM/370 Programming Functions

#### New Programming Features

- The Virtual = Real Performance option
- The Dedicated Channel Performance option
- The Virtual and Real Channel-to-Channel Adapter

### Support for the Following Devices

#### New Hardware Features

- The IBM 3211 Printer
- The IBM 3410/3411 Magnetic Tape Subsystem
- The IBM System/370 Models 155 II and 158

### Recovery Management Support

#### Maintenance: Program and Documentation

The section on Recovery Management Support has been rewritten to include the following changes and additions:

- A revised explanation of the initial state of the recovery mode for main storage errors.
- A revised termination procedure where recovery via an automatic restart is attempted before placing the system in a disabled wait state.
- The addition of a Buffer Error Routine as part of the Machine Check Handler to perform error recovery on those CPUs that have high speed buffers.

### Additional Modules

#### New: Program and Documentation

The following modules have been added as part of the Control Program:

- The new module, DMKCPB, now simulates the SYSTEM, EXTERNAL, READY, NOTREADY, RESET, and REWIND commands for the virtual machine. This function has been removed from module DMKCFM.
- The new module, DMKTRC, now contains all the TRACE processing routines. The initialization, modification, and termination of the TRACE function remains in module DMKTRA.
- The new module, DMKCFG, contains the SAVESYS command processing routine. This function was removed from module DMKCPV.

### Additional Data Area

#### New: Program and Documentation

The "Extended Outboard Recording Block" (XOBR3211) has been added as a continuation of the "I/O Error Information Block" (IOERBLOK). It will hold additional sense data for devices that return more than 24 sense bytes.

### Additional and Revised Flowcharts

#### Maintenance: Program and Documentation

The flowcharts and module/entry point directory entries for the following modules have been revised to reflect the above-cited new features and support:

DMKCCH	DMKDSP	DMKSPL
DMKCCW	DMKFRE	DMKTAP
DMKCDS	DMKGEN	DMKTDK
DMKCFG	DMKIOE	DMKTRA
DMKCFM	DMKIOF	DMKTRC
DMKCFP	DMKIOG	DMKUSO
DMKCNS	DMKIOS	DMKVAT
DMKCPB	DMKLDG	DMKVCA
DMKCPV	DMKMCH	DMKVCH
DMKCPV	DMKNEM	DMKVCN
DMKCSO	DMKPAG	DMKVDB
DMKDas	DMKPRG	DMKVDS
DMKDEF	DMKPRV	DMKVIO
DMKDGd	DMKPSA	DMKVMI
DMKDIA	DMKRSP	DMKVSP
DMKDMP	DMKSCN	

### Error Messages and Codes

#### Maintenance: Program and Documentation

The following Error Messages have been added:

DMKCCH605I	DMKDas956A
DMKCCH606I	DMKDIA011E
DMKCFG044E	DMKDMP909W
DMKCFG170E	DMKMCH610I
DMKCFG171E	DMKMCH611I
DMKCFG172E	DMKMCH612W
DMKCFG173E	DMKMCH614I
DMKCFG435E	DMKMCH616I
DMKCFP174E	DMKPAG415E
DMKCPB005E	DMKPRG453W
DMKCPB006E	DMKSPL501I
DMKCPB012E	DMKSPL503A
DMKCPB022E	DMKSPL504A
DMKCPB026E	DMKSPL529I
DMKCPB040E	DMKUDR475I
DMKCPV144W	DMKVDB034E
DMKCSO036E	DMKWRRM911W

(See Over)

The following Error Messages have been deleted:

DMKCFP005E	DMKDIA110E
DMKCFP006E	DMKMCH610W
DMKCFP012E	DMKMCH611W
DMKCFP022E	DMKMCH612I
DMKCPV044E	DMKMCH614W
DMKCPV170E	DMKMCH616W
DMKCPV171E	DMKMCH620I
DMKCPV172E	DMKSPL517I
DMKCPV173E	DMKWRM910W

The following Wait state codes have been added:

00D            00F

The following ABEND codes have been added:

BLD001	DSP004	PTR010
CFM001	FRE010	SCH001
CNS008	PTR008	TRC001
DSP003	PTR009	

The following ABEND code has been deleted:

TRA001

#### Miscellaneous

Maintenance: Documentation Only

This edition includes other minor technical and typographical changes too numerous to list.

#### NEW DEVICE SUPPORT

##### New: Programming Feature

The IBM System/370 Model 168; the IBM 2860, 2870, 2880 standalone channels; and the IBM 2305 Fixed Head Storage, Model 1, are now supported.

#### USER ACCOUNTING OPTION

##### New: Programming Feature

It is now possible for a user to charge another user for CPU time. A new diagnose code (4C) is provided for this function. This option is described under "Accounting Card Processing" in the section on the "Real Spooling Manager." The new diagnose code is described in the section "Privileged Instructions."

#### VIRTUAL CONSOLE SPOOLING

##### New: Programming Feature

The virtual console is now supported for spooling operations. Documentation of this support appears in the section, "Virtual Spooling Manager."

#### CP INTERNAL TRACE TABLE IMPLEMENTATION

##### New: Programming Feature

A new CP command, MONITOR, allows the user to stop and restart the recording of real machine events in the internal trace table. Previously, the tracing was always active. This feature is described in the section, "CP Internal Trace Table."

#### STOP OPERAND AND PARAMETER PASSING FOR THE CP IPL COMMAND

##### New: Programming Feature

The STOP operand in the CP IPL command will halt execution and allow parameters to be passed resulting in the loading of an alternate nucleus. This function is described in "LOGON of User" in the section "System User Interface."

#### PERFORMANCE ENHANCEMENT

##### Maintenance: Program and Documentation

The Dispatcher/Scheduler routines have been modified to improve performance. The section on the Dispatcher/Scheduler has been rewritten.

#### ADDITIONAL MODULES

##### New: Program and Documentation

The following modules have been added as part of the Control Program:

- The modules DMKSIX, DMKSEV, and DMKEIG handle the channel logout analysis for the 2860, 2870, and 2880 standalone channels, respectively.
- The module, DMKGRA, handles VM/370 console spooling.
- The module, DMKLOC, locks and unlocks a system resource. This code was previously restricted to use in DMKUDR.
- The module, DMKMCC, handles the new CP command,

MONITOR.

- The module, DMKRSE, retries and attempts recovery for real U/R device I/O errors. This function was originally in module DMKSPL.

ADDITIONAL DATA AREAS

New: Program and Documentation

ACCTBLOK -- User Accounting Block

New: Documentation Only

CCHREC -- Channel Check Handler Record

MCHAREA -- Machine Check Save Area

MCRECORD -- Machine Check Handler Record

The above blocks are defined in the "Data Areas -- Control Blocks" section.

ADDITIONAL AND REVISED FLOWCHARTS

New: Program and Documentation

DMKEIG DMKMCC DMKSIX

DMKGRA DMKRSE

DMKLOC DMKSEV

Maintenance: Program and Documentation

DMKACO DMKDSO DMKSPL

DMKCCH DMKHVC DMKTMR

DMKCDS DMKIOE DMKUDR

DMKCFP DMKIOG DMKUSO

DMSCFS DMKIOS DMKVCH

DMKCKP DMKLNK DMKVCN

DMKCPB DMKLOG DMKVDB

DMKCPI DMKMCH DMKVDS

DMKCPV DMKPRG DMKVM I

DMKCSP DMKRSP DMKVSP

DMKDEF DMKSCH DMKW RM

ERROR MESSAGES AND CODES

Maintenance: Program and Documentation

The following error messages have been added:

DMKCFP177E	DMKRSE503I
DMKCPB059E	DMKRSE504A
DMKMCC002E	DMKRSE504I
DMKMCC026E	DMKRSE505A
DMKMCH003E	DMKRSE508I
DMKMCH026E	DMKRSE520A
DMKRSE500I	DMKRSE520I
DMKRSE501A	DMKRSE521I
DMKRSE501I	DMKRSE524I
DMKRSE502I	DMKRSE525I
DMKRSE503A	DMKRSE529I

The following error messages have been deleted:

DMKSPL500I	DMKSPL505D
DMKSPL501A	DMKSPL508I
DMKSPL501I	DMKSPL520I
DMKSPL502D	DMKSPL521I
DMKSPL503A	DMKSPL524I
DMKSPL503I	DMKSPL525I
DMKSPL504A	DMKSPL529I
DMKSPL504I	

The following ABEND codes have been added:

IOS001	IOS003
IOS002	UDR001

MISCELLANEOUS

Maintenance: Documentation Only

This Technical Newsletter contains other minor technical and typographical changes, too numerous to mention.



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## INTRODUCTION

The VM/370 Control Program (CP) manages the resources of a System/370 in order to provide virtual storage support through the implementation of virtual machines. This support is implemented in such a way that each terminal user appears to have the complete functional capabilities of a dedicated System/370 at his disposal, even though many other users may be running batch, teleprocessing, time sharing testing, or production jobs at the same time.

A user defines the configuration he requires -- input/output (I/O) device addresses, and a storage size up to 16 million bytes -- regardless of whether they match the real machine's configuration. Virtual devices must have real counterparts, but not always in a one-for-one ratio. For example, many users' readers, punches, and printers can be mapped onto common spool disks, and their virtual disk devices may be mapped as minidisks onto different sections of common disk packs, effectively multiplying the number of logical disk devices that are available on the real machine.

Each user's virtual machine comprises

- An operator's console (his remote terminal)
- A virtual CPU either with or without the Virtual Storage Addressing feature
- Virtual storage of up to 16 million bytes
- Virtual I/O devices

Virtual I/O devices are controlled by the virtual machine's operating system, not by the VM/370 control program. Thus, the support for the proper number and type of I/O devices must be provided by the operating system of the virtual machine for proper operation. The VM/370 control program monitors, translates, and schedules all real I/O operations to provide system integrity. It executes all virtual machine operation in a problem state by trapping, screening, and processing all the interrupts, and passing on the necessary information to the appropriate virtual machine. Only

the VM/370 control program executes in the privileged state.

In order to increase the amount of real main storage available to user programs, parts of the VM/370 control program that are infrequently used are not required to be resident in main storage. Instead, they reside on part of the paging auxiliary storage used by the system, and are brought into main storage only when their functions are required.

Since the VM/370 Control Program nonresident modules are effectively paged into main storage, the control program itself must have virtual storage space associated with it. This space is anchored at the System VMBLOK, which is assembled into the resident control program in the module DMKSYS. The VMBLOK has a pointer to a segment table, which in turn references a set of page and swap tables that describe CP's virtual storage space.

The virtual space is divided into 2 parts; the first 4 segments (256K) is reserved for executable control program code, both resident and pageable; the remaining storage (at least another 256K) is dynamically allocated for spooling buffers and for user directory functions. In order for a routine to be pageable, a number of restrictions must be observed.

When the system is loaded, resolved, and written onto the system residence volume, those modules that are to be pageable must be loaded at addresses higher in main storage than the symbol DMKCPEND, which defines the last byte of the resident CP nucleus. This arrangement can be accomplished by reordering the LOADLIST EXEC used by the VMFLOAD procedure when punching out the text decks that will compose the CP system. Any pageable modules are listed after the entry for DMKCPE. In addition, each pageable module must be preceded by the 'SPB' loader control card. This 'Set Page Boundary' card forces the loader to start loading the succeeding module at the next higher 4k page boundary and ensures that the entire module will be resident when it is paged in.

If several pageable modules perform similar or related functions and it is felt that they are likely to be resident at the same time, they may be included in the same page by omitting the SPB cards that would normally have preceded the 2nd and subsequent modules. The group of modules to be loaded together must not exceed 4K as their total storage requirement; if they do, one or more must be loaded in separate pages, since no page boundary crossover in pageable control program is allowed. All currently pageable CP modules punch their own SPB card via an assembler PUNCH statement, except those that are designed to reside in a page along with other modules.

#### CP INITIALIZATION

The function of system initialization (IPL) is to prepare VM/370 for operation. Some of the tasks to be performed are:

- Main storage must be initialized
- Devices must be mounted
- Warm start records must be read from the warm start cylinder
- Space must be allocated for the system dump file
- The system operator must be logged on

In the case of a system restart following a failure, active files and the system log message must be written to the checkpoint cylinder before the Control Program nucleus can be brought into main storage. The user can now logon.

#### VIRTUAL MACHINE CONTROL

A virtual machine is created for a user when he logs into VM/370, on the basis of information stored in his user directory entry. The entry for each user identification includes a list of the virtual I/O devices associated with the particular virtual machine and the real device mappings.

Additional information about the virtual machine is maintained in the directory file. Included are the VM/370 command privilege class, accounting data, normal and maximum virtual storage sizes, and optional virtual machine characteristics such as extended control mode.

The Control Program supervises the execution of virtual machines by (1) permitting only problem state execution except in its own routines, and (2) receiving control after all real computing system interrupts. CP intercepts each privileged instruction and simulates it if the current program status word of the issuing virtual machine indicates a virtual supervisor state; if the virtual machine is executing in virtual problem state, the attempt to execute the privileged instruction is reflected back to the virtual machine as a program interrupt. All virtual machine interrupts (including those caused by attempting privileged instructions) are first handled by CP, and are reflected to the virtual machine if an analogous interrupt would have occurred on a real machine.

#### Virtual Machine Time Management

The real CPU is time sliced to simulate multiple virtual CPUs. Virtual machines that are executing in a conversational manner are given access to the real CPU more frequently than those that are not; these conversational machines are assigned the smaller of two possible time slices. CP determines execution characteristics of a virtual machine at the end of each time slice on the basis of the recent frequency of its console requests or terminal interrupts. The virtual machine is queued for subsequent CPU utilization according to whether it is a conversational or nonconversational user of system resources.

A virtual machine can gain control of the CPU only if it is not waiting for some activity or resource. The virtual machine itself may enter a virtual wait state after an I/O operation has begun. The virtual machine cannot gain control of the real CPU if it is waiting for a page of storage, if it is waiting for an I/O operation to be translated and started, or if it is waiting for a CP command to finish execution.



A virtual machine can be assigned a priority of execution. Priority is a parameter affecting the execution of a particular virtual machine as compared with other virtual machines that have the same general execution characteristics. Priority may be assigned by the real machine operator, but is more frequently a parameter of the virtual machine's directory entry.

### Virtual Machine Storage Management

The normal and maximum storage sizes of a virtual machine are defined as part of the virtual machine configuration in the VM/370 directory. The virtual storage size can be temporarily redefined to any value that is a multiple of 4K and not greater than the maximum defined value. VM/370 implements this storage as virtual storage. The storage may appear as paged or nonpaged to the virtual machine, depending upon whether the extended control mode option has been specified for that virtual machine. This option is required if operating systems that control virtual storage, such as OS/VS1 or VM/370, are to be run in the virtual machine.

Storage in the virtual machine is logically divided into 4096 byte areas called pages. A complete set of segment and page tables is used to describe the storage of each virtual machine. These tables are maintained by CP and reflect the allocation of virtual storage pages to blocks of real storage. Virtual storage addressing is accomplished through use of these tables by the System/370 machine. Storage in the real machine is logically and physically divided into 4096 byte areas called page frames or blocks.

Only referenced virtual storage pages are kept in real storage, thus optimizing real storage utilization. Further, a page can be brought into any available page frame; the necessary relocation is done during program execution by a combination of VM/370 and dynamic address translation on the System/370. The active pages from all logged-in virtual machines and from the pageable routines of CP compete for available page frames. When the number of page frames available for allocation falls below a threshold value, CP determines which virtual storage pages currently allocated to real storage are relatively inactive and initiates suitable page-out operations for them.

Inactive pages are maintained on a direct access storage device. If an inactive page has been changed at some time during virtual machine execution, CP assigns it to a paging device, selecting the fastest such device with available space. If the page has not changed, it remains allocated in its original direct access location and is paged into real storage from there the next time the virtual machine references that page. A virtual machine program can use the DIAGNOSE instruction to communicate to CP that the information from specific pages of virtual storage is no longer needed; CP then releases the areas of the paging devices which had been assigned to hold the specified pages.

Paging is done on demand by CP. This means that a page of virtual storage is not read (paged) from the paging device to a real storage block until it is actually needed for virtual machine execution. No attempt is made by CP to anticipate what pages might be required by a virtual machine. While a paging operation is being performed for one virtual machine, another virtual machine can be executing. Any paging operation initiated by CP is transparent to the virtual machine.

If the virtual machine is executing in extended control mode with translate on, then two additional sets of segment and page tables are maintained. The virtual machine operating system is responsible for mapping the virtual storage created by it to the storage of the virtual machine. CP uses this set of tables in conjunction with the page and segment tables created for the virtual machine at login time to build shadow page tables for the virtual machine. These shadow tables map the virtual storage created by the virtual machine operating system to the storage of the real computing system. The tables created by the virtual machine operating system may describe any page and segment size permissible in the IBM System/370.

The system operator may assign the reserved page frames option to a single virtual machine. This option, specified by the SET RESERVE command, assigns a specific amount of the storage of the real machine to the virtual machine. CP dynamically builds a set of reserved real storage page frames for this virtual machine during its execution until the maximum number "reserved" has been reached. Since other virtual machines' pages are not allocated from this reserved

set, the effect is that the most active pages of the selected virtual machine remains in real storage.

During the process of CP system generation, the installation may specify that a single virtual machine is to be given an option called virtual=real. With this option, the virtual machine's storage is allocated directly from real storage at the time CP is initially loaded, and remains so allocated unless released via operator command. All pages except page zero are allocated to the corresponding real storage locations. In order to control the real computing system, real page zero must be controlled by CP. Consequently, the real storage size must be large enough to accommodate the CP nucleus, the entire Virtual=Real virtual machine, and the remaining pageable storage requirements of CP and the other virtual machines.

The virtual=real option improves performance in the selected virtual machine since it removes the need for CP to perform paging operations for the selected virtual machine. The virtual=real option is necessary whenever programs that contain dynamically modified channel programs (excepting those of OS ISAM) are to execute under control of CP.

#### Virtual Machine I/O Management

A real disk device can be shared among multiple virtual machines. Virtual device sharing is specified in the directory entry or by a user command. If specified by the user an appropriate password must be supplied before gaining access to the virtual device. A particular virtual machine may be assigned read-only or read/write access to a shared disk device. CP verifies each virtual machine I/O operation against the parameters in the virtual machine configuration to ensure device integrity.

The virtual machine operating system is responsible for the operation of all virtual devices associated with it. These virtual devices may be defined in the directory entry of the virtual machine, or they may be attached to (or detached from) the virtual machine's configuration while it remains logged on. Virtual devices may be dedicated, as when mapped to a fully

equivalent real device; shared, as when mapped to a minidisk or when specified as a shared virtual device; or spooled by CP to intermediate direct access storage.

In a real machine running under control of OS, I/O operations are normally initiated when a problem program requests OS to issue a START I/O instruction to a specific device. Device error recovery is handled by the operating system. In a virtual machine, OS can perform these same functions, but the device address specified and the storage locations referenced are both virtual. It is the responsibility of CP to translate the virtual specifications to real.

In addition, the interrupts caused by the I/O operation are reflected to the virtual machine for its interpretation and processing. If I/O errors occur, CP records them but does not initiate error recovery operations. These are the responsibility of the virtual machine operating system.

I/O operations initiated by CP for its own purposes (paging and spooling), are performed directly and are not subject to translation.

#### Spooling

A virtual unit record device, which is mapped directly to a real unit record device, is said to be dedicated. The real device is then controlled completely by the virtual machine's operating system.

CP facilities allow multiple virtual machines to share unit record devices. Since virtual machines controlled by CMS ordinarily have modest requirements for unit record I/O, such device sharing is quite advantageous, and it is the standard mode of system operation.

Spooling operations cease if the direct access storage space assigned to spooling has been exhausted, and the virtual unit record devices appear in a not ready status. The system operator may make additional spooling space available by purging existing spool files or by assigning additional direct access storage space to the spooling function.

Specific files can be transferred from the spooled card punch or printer of a virtual machine to the card reader of the same or another virtual machine. Files transferred between virtual unit record devices by the spooling routines are not physically punched or printed. With this method, files can be made available to multiple virtual machines, or to different operating systems executing at different times in the same virtual machine.

CP spooling includes many desirable options for the virtual machine user and the real machine operator. These options include printing multiple copies of a single spool file, backspacing any number of printer pages, and defining spooling classes for the scheduling of real output.

### Console Functions

The CP console functions allow the user to control the virtual machine from the terminal, much as an operator controls a real machine. Virtual machine execution can be stopped at any time by use of the terminal's attention key; it can be restarted by typing in the appropriate CP command. External, attention, and device ready interrupts can be simulated on the virtual machine. Virtual storage and virtual machine registers can be inspected and modified, as can status words such as the PSW and the CSW. Extensive trace facilities are provided for the virtual machine, as well as a single-instruction mode. Commands are available to invoke the spooling and disk sharing functions of CP.

Console functions are divided into privilege classes. The directory entry for each user assigns one or more privilege classes. The classes are:

- System operator
- Operator
- System programmer
- Spooling operator
- Systems analysts

- Customer engineering
- General users

Commands in the system analysts class may be used to inspect real storage locations, but may not be used to make modifications to real storage. Commands in the operator class provide real resource control capabilities. System operator commands include all those relating to virtual machine performance options, such as assigning a set of reserved page frames to a selected virtual machine. See the "CP Commands" sections of this chapter for more information.

### PROGRAM STATES

When instructions in the Control Program are being executed, the real computer is in the supervisor state; at all other times, when running virtual machines, it is in the problem state. Therefore, privileged instructions can only be executed by the Control Program. Programs running on a virtual computer can issue privileged instructions; such an instruction causes an interruption that is handled by the Control Program. CP examines the operating status of the virtual machine PSW. If the virtual machine indicates that it is functioning in supervisor mode, then the privileged instruction is simulated according to its type. If the virtual machine is in problem mode, then the privileged interrupt is reflected to the virtual machine.

Only the Control Program may operate in the supervisor state on the real machine. All programs other than CP operate in the problem state on the real machine. All user interrupts, including those caused by attempted privileged operations, are handled by CP, which then reflects to the user program only those interrupts that the user program would expect from a real machine. A problem program executes on the virtual machine in a manner identical to its execution on a real System/370 CPU, as long as it does not violate the CP restrictions.

PREFERRED VIRTUAL MACHINE

CP supports four special virtual machine operating environment functions. Each function can be applied to one virtual machine at a time. Although each function could be applied to a different virtual machine, optimum performance would not be achieved. Each function is discussed separately following.

## FAVORED EXECUTION

CP attempts to provide a specified percentage of CPU time to a particular virtual machine. CP attempts to provide up to the specified percentage of CPU time to a particular virtual machine, provided that the virtual machine is functioning so that it can fully utilize the CPU time. At regular time intervals the CP dispatcher checks the CPU time used by the particular virtual machine. If the specified percentage is exceeded, the machine becomes the lowest priority user in the system. If the percentage used is lower than that specified, the virtual machine has highest priority execution for the remainder of the interval. The percentage of CPU time assured is specified in the privileged class command that invokes the function.

CP can also assure that a designated user will never be dropped from the active (in queue) subset by the scheduler. When the user is runnable, he is placed in the dispatchable list at his normal priority.

## RESERVED PAGE FRAMES

CP uses chained lists of table entries for available and pageable pages. Pages for users are assigned from the available lists which is replenished from the pageable list.

Pages which are temporarily locked in real storage are not available or pageable. Paging proceeds using demand paging with a "reference bit" algorithm to select the best page for swapping. The reserved page frames option gives a particular virtual machine an

essentially "private" set of pages. The pages are not locked, that is, they can be swapped, but usually only for the specified virtual machine. The number of reserved pages for the virtual machine are specified as a maximum. The page selection routine will select an available page for a reserved user and mark that page "reserved" if the maximum specified for the user has not been reached. If an available, unreferenced "reserved" page is encountered during page replenishment for the reserved user, it is used whether or not the maximum has been reached. If the page selection routine cannot locate an available page for other users because they are all "reserved", the routine may have to steal the reserved pages.

## DEDICATED CHANNELS

Since the devices on a channel are often shared between virtual machines (minidisks and dedicated devices) and shared with system functions (paging and spooling), CP schedules all the I/O requests to achieve a balance between machines. In addition, CP simulates the reflection of the subsequent I/O interrupts to the virtual machines. By specifying a dedicated channel(s) for a virtual machine, the CP channel scheduling function is bypassed. The virtual device addresses on the dedicated channel must match the real device addresses. Since the channels are dedicated, CP uses the virtual machine masking to control the real channel masking. I/O interrupts from the dedicated channel are presented in the order of occurrence using a single element stack and the real channel masking.

A single virtual machine may have multiple dedicated channels. Also, multiple virtual machines may each have a separate dedicated channel.

## VIRTUAL=REAL

This feature requires that the CP nucleus be reorganized to provide a "hole" in real storage large enough to contain the entire storage area of the virtual machine. For the virtual machine, each page from page 1 to the last page (n) is in its true real

storage location; only page zero is relocated. The virtual machine is still run in relocate mode, but since the virtual page address is the same as the real page address, no CCW translation is required for the virtual machine. Since no CCW translation is performed, no check is made of the I/O data addresses. The virtual machine must ensure that no I/O data transfer is specified into page zero or into any page not in the virtual machine's domain.

There are several considerations for the virtual=real option of preferred machine support that affect overall system operation:

- The area of contiguous storage built for the virtual=real machine must be large enough to contain the entire addressing space of that machine.
- While allocated as such, the storage reserved for the virtual=real machine can only be used by a virtual machine with that option. It is not available to other users for paging space nor for VM/370 usage, even when the virtual=real machine is not logged on. For this reason, it is expected that the virtual=real machine will be a high availability, high throughput machine.

The virtual=real storage can be released by the operator. That storage is then available for paging. Once virtual=real storage space is released by the operator, a VM/370 IPL is necessary to again allocate that storage to that virtual=real machine.

- The virtual machine with the virtual=real option operates in the pre-allocated storage area with normal CCW translation in effect until the execution of the SET NOTRANS ON command. At that time, all subsequent I/O operations are performed from the virtual CCWs in the virtual=real space without translation. In this mode, the virtual machine must not perform I/O operations into page zero nor beyond its addressable limit. Violation of this requirement causes destruction of the VM/370 system and/or other virtual machines.
- If the virtual=real machine performs a virtual reset or IPL, then the normal CCW translation is performed until the issuance of the SET NOTRANS ON command.

Only the virtual=real virtual machine can issue the command. A message is issued if normal translation mode is entered.

#### CP INTERRUPTION HANDLING

##### I/O INTERRUPT

I/O interrupts from completed I/O operations initiate various completion routines and the scheduling of further I/O requests. The I/O interrupt handling routine also gathers device sense information.

##### PROGRAM INTERRUPT

Program interrupts can occur in two states. If the CPU is in supervisor state, the interrupt indicates a system failure in the CP nucleus and causes a system abend. If the CPU is in problem state, then a virtual machine is executing. If the program interrupt indicates that the Dynamic Address Translation (DAT) feature has an exception, a virtual machine issued a privileged instruction, or a protection exception occurred for a shared segment system, then CP takes control to perform any required processing to satisfy the exception. Usually, the interrupt is transparent to the virtual machine execution. Most other program interrupts result from virtual machine processing and are reflected to the machine for handling. For a complete discussion of this subject, see the appropriate explanation in the section "Method of Operation".

##### MACHINE CHECK INTERRUPT

When a machine check occurs, the CP Recovery Management Support (RMS) gains control to save data associated with the failure for FE maintenance. RMS analyzes the failure to determine the extent of damage.

Damage assessment results in one of the following actions being taken:

- System Termination
- Selective Virtual User Termination
- Refreshing of damaged information with no affect on system configuration
- Refreshing of damaged information with the defective storage page removed from further systems use
- Error recording only for certain soft machine checks

The system operator is informed of all actions taken by the RMS routines. When a machine check occurs during VM/370 startup (before the system is set up well enough to permit RMS to operate successfully), the CPU goes into a disabled wait state and places a completion code of X'00B' in the high-order bytes of the current PSW.

#### SVC INTERRUPT

When an SVC interrupt occurs, the SVC interrupt routine is entered. If the machine is in problem mode, the type of interrupt is reflected back to the pseudo-supervisor (that is, the supervisor operating in the user's virtual machine). If the machine is in supervisor mode, the SVC interrupt code is determined, and a branch is taken to the appropriate SVC interrupt handler.

#### EXTERNAL INTERRUPT

If a timer interrupt occurs, CP processes it according to type. The interval timer indicates time-slice end for the running user. The clock comparator indicates that a specified timer event has occurred, such as midnight, scheduled shutdown, or user event reached. The CPU timer indicates that a virtual machine's allowed execution interval (time in queue) has expired.

The external console interrupt invokes CP processing to switch from the 3210 or 3215 to an alternate operator's console.

#### FREE STORAGE MANAGEMENT

During its execution, CP occasionally requires small blocks of storage that are used for the duration of a task. This storage is obtained from the free storage area. The free storage area is divided into various size subpools. The requestor informs the Free Storage Manager the size of the block required and the smallest available subpool that fulfills the request is allocated to the requestor. When the block is no longer needed, the requestor informs the Free Storage Manager and the block is returned to free storage.

If the request for free storage cannot be fulfilled the Free Storage Manager requests the temporary use of a page of storage from the Dynamic Paging Area. If a page is obtained, then the page is chained to the free storage area and used for that purpose until it is no longer needed and subsequently returned to the Dynamic Paging Area.

If the request for a page cannot be fulfilled, the requestor waits until free storage becomes available.

#### EXECUTING THE PAGEABLE CONTROL PROGRAM

Calls to pageable routines are recognized at execution time by the SVC 8 linkage manager in DMKPSA. For every SVC 8, the called address (in the caller's GPR15) is tested to see if it is within the resident nucleus. If it is less than DMKCPEND and greater than DMKSLC, the called routine's base address is placed in GPR12 and control is passed to the called routine in the normal way. However, if the called address is above DMKCPEND or below DMKSLC, the linkage manager issues a TRANS macro, requesting the paging manager to locate and, if necessary, page-in the called routine. The TRANS is issued with LOCK option. Thus, the lock count associated with the called routine's real page indicates the responsibility count of the module.

- When the module is called, the count is incremented.
- When the routine exits via SVC 12, the count is decremented.

When the count reaches zero, the pageable routine is unlocked and is eligible to be paged out of the system. However, since all CP pageable modules are reentrant, the page is never swapped out, but when stolen is placed directly on the free page list.

Since unlocked pageable routines participate in the paging process in a manner similar to user virtual storage pages, the Least Recently Used approximation used by page selection tends to make highly used control program routines, even when not locked, remain resident. The called routine is locked into real storage until it exits. Thus, it can request asynchronously scheduled function, such as I/O or Timer interrupts, as long as it dynamically establishes the interrupt return address for the requested operation and does not give up control via an EXIT macro prior to receiving the requested interrupt.

Addressability for the module while it is executing is guaranteed since the CALL linkage loads the real address of the paged module into GPR12 (the module base register) prior to passing control. If all addressing is done in a base/displacement form, the fact that the module is executing at an address different from that at which it was loaded is transparent. Although part of the control program is pageable it never runs in relocate mode. Thus, the CPU is not degraded by the DAT feature being active, and there is no problem of handling disabled page-faults.

#### SYSTEM SUPPORT MODULES

The system support modules provide CP several common functions in the area of data conversion and control block scanning and verification. Since most of the routines operate at the lowest level of control, they are linked to via the BALR option of the CALL macro, and make use of the BALRSAVE and TEMPSAVE workareas in DMKPSA. Two exceptions to this are the virtual and real I/O control block scan routines DMKSCNVU and DMKSCNRU. These routines do not alter the contents of the BALRSAVE area, and hence may be called by another low level BALR routine.

#### CONTROL REGISTER USAGE

Every IBM System/370 CPU provides the program with 16 logical control registers (logical registers since the number that are active depends on the features installed in the machine at any one time) that are addressable for loading and storing from BC mode. VM/370 provides only a single control register, control register zero, for normal virtual machines, for processing systems that do not require the full set of registers (for example, CMS, DOS, or other operating systems for System/360).

Any user whose virtual machine operating system requires the use of control registers other than control register zero can request the full set of 16 registers by specifying the ECMODE option in the VM/370 user-directory entry for his virtual machine. Specifying this option does not imply that the virtual machine will encounter any of the additional overhead associated with use of the Extended Control mode but permits the use of all 16 control registers from either BC or EC mode.

A virtual machine, which utilizes any System/370 features that use the control registers, requires the ECMODE option. Some of these features are expanded timer support of the System/370, (CPU timer, clock comparator, etc.), the virtual relocate-mode and its instructions, RRB, LRA, PTLB, virtual monitor calls, virtual Program Event Recording (PER), etc.

#### RESTRICTIONS AND CONVENTIONS FOR PAGEABLE CP MODULES

CP modules that are to be pageable must observe the following restrictions and conventions when they are designed and coded:

1. The module should be completely reentrant. Any messages to be modified, temporary work or scratch areas, or program switches must be allocated from system free storage or from the caller's save area.

2. The module must be entered via the standard SVC 8 CALL linkage. Modules entered via BALR or GOTO cannot be pageable.
3. The module cannot contain any A or V type address constants that point to locations within itself or within other pageable modules, and it cannot contain any CCWs that contain data address within itself. The only exceptions are address constant literals generated as the result of CALLs to other modules (since these addresses are dynamically relocated at execution time, they must be resolved by the loader to the loaded address of the called module) and a pageable module that locks itself into storage. In practice, this restriction simply means that data or instructions within the pageable routine must be referenced via base/displacement addressing, and the address in register 15 for a CALL may not be generated via a LOAD ADDRESS.
4. The pageable module must be no more than 4096 bytes in length.

If the above design and coding restrictions are adhered to, the CP module can be added to the existing pageable nucleus modules by utilizing the service routine, VMFLOAD, which is described in the "VM/370 Maintenance Procedures" chapter of the publication IBM Virtual Machine Facility/370: Service Routines Program Logic, Order No. SY20-0882. Additional information can be found in "Appendix I" of the publication IBM Virtual Machine Facility/370: Planning and System Generation Guide, Order No. GC20-1801.

## MODULES

### Executable Resident Modules

DMKCCH  
 DMKCCW  
 DMKCFM  
 DMKCNS  
 DMKCVT  
 DMKDAS  
 DMKDGD  
 DMKDMP  
 DMKDSP  
 DMKFRE  
 DMKGEN  
 DMKHVC  
 DMKIOE  
 DMKIOS  
 DMKMCH  
 DMKMSW  
 DMKPAG  
 DMKPGS  
 DMKPGT  
 DMKPRG  
 DMKPRV  
 DMKPSA  
 DMKPTR  
 DMKQCN  
 DMKRPA  
 DMKRSP  
 DMKSCH  
 DMKSCN  
 DMKSTK  
 DMKTMR  
 DMKUNT  
 DMKVAT  
 DMKVCN  
 DMKVIO  
 DMKVSP



### Executable Pageable Modules

DMKACO  
DMKBLD  
DMKCDB  
DMKCDS  
DMKCFD  
DMKCFG  
DMKCFP  
DMKCFS  
DMKCFT  
DMKCKP  
DMKCPB  
DMKCPI  
DMKCPV  
DMKCQG  
DMKCQP  
DMKCSO  
DMKCSP  
DMKCSU  
DMKDEF  
DMKDIA  
DMKDRD  
DMKEIG  
DMKERM  
DMKGRA  
DMKIOF  
DMKIOG  
DMKISM  
DMKLNK  
DMKLLOC  
DMKLOG  
DMKMCC  
DMKMID  
DMKMSG  
DMKNEM  
DMKRSE  
DMKSAV  
DMKSEP  
DMKSEV  
DMKSIX  
DMKSPL  
DMKTAP  
DMKTDK  
DMKTRA  
DMKTRC  
DMKTRM  
DMKUDR  
DMKUSO  
DMKVCA

DMKVCH  
DMKVDB  
DMKVDS  
DMKVMI  
DMKWRM

### Data Area Modules

In addition to the executable resident and pageable modules there are certain modules that only contain data areas and do not execute. These modules are:

#### Resident

##### Module    Contents

DMKCPE    Defines the end of the CP nucleus  
DMKRIO    I/O device blocks  
DMKSYS    System constants  
DMKTBL    Terminal translate table

#### Pageable

##### Module    Contents

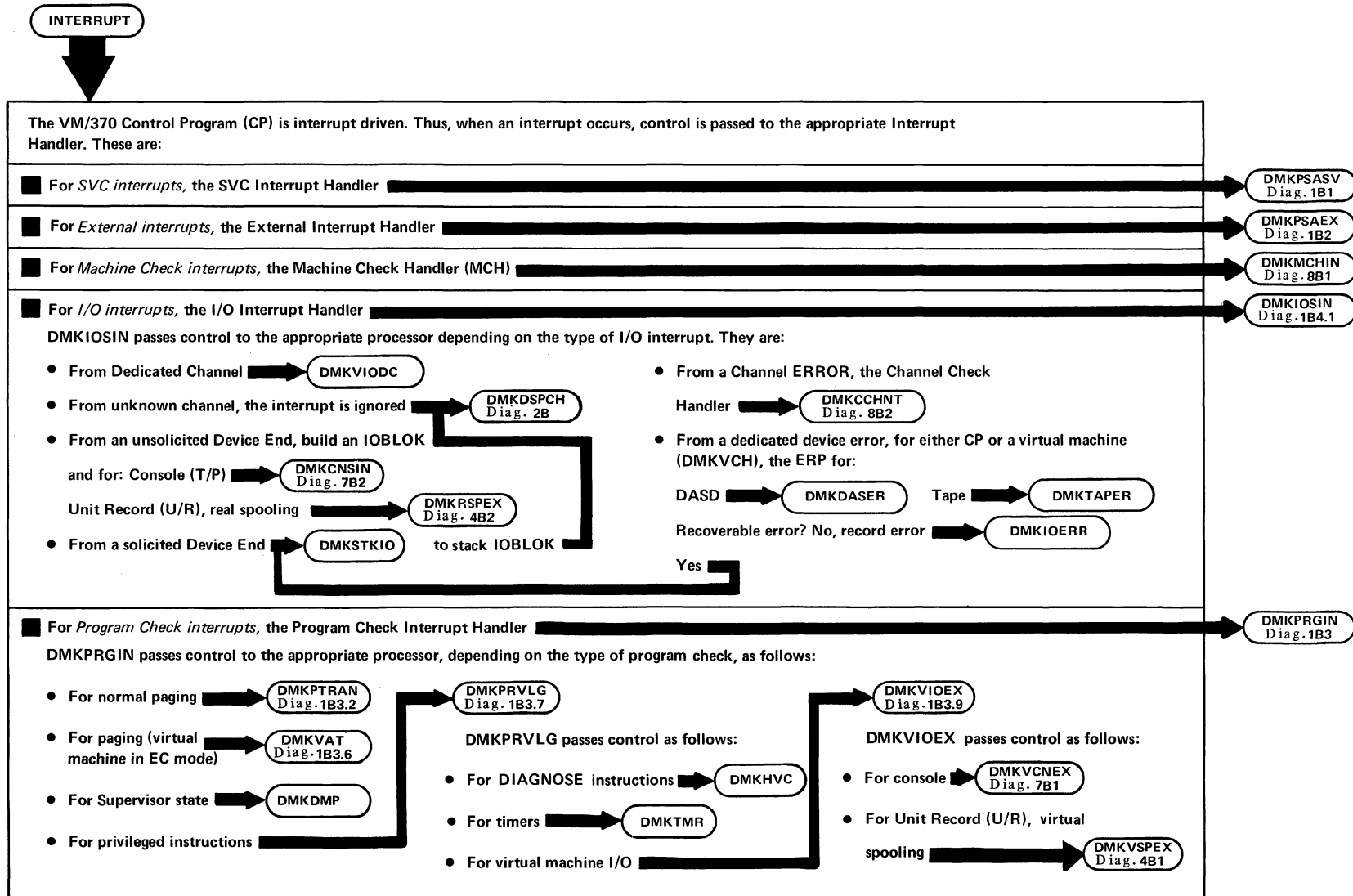
DMKBOX    Output separator table  
DMKFCB    3211 Forms control Buffer (FCB) load tables  
DMKSNT    System name table  
DMKSYM    System symbol table  
DMKUCB    3211 Universal Character Set Buffer (UCSB)  
          load tables  
DMKUCS    1403 Universal Character Set (UCS) load  
          tables

The data areas within these modules are discussed throughout this publication; most are illustrated in the section "Data Areas".

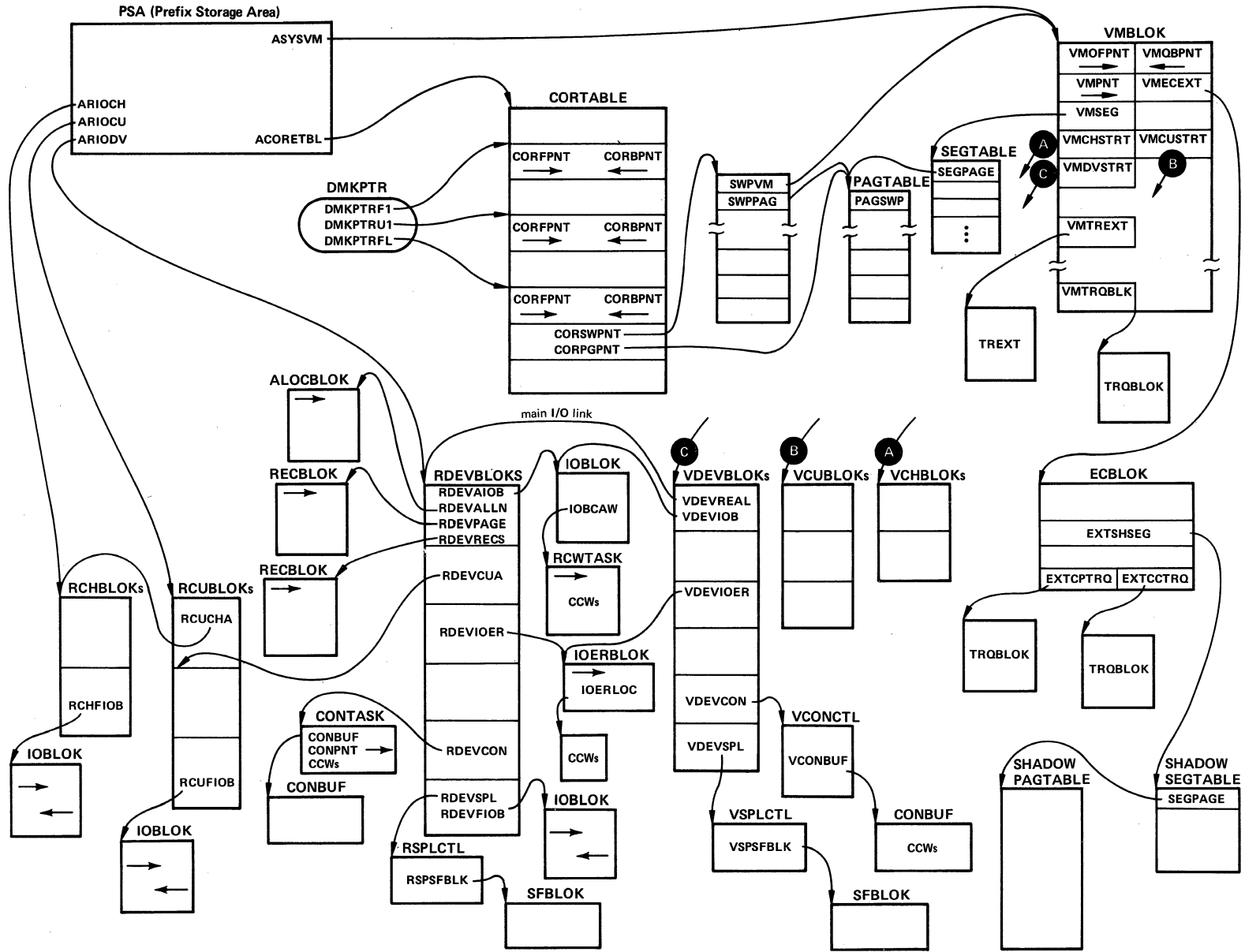


**METHOD OF OPERATION**

**Diag. 1A. Overview of Method of Operation Diagrams**



Diag. 1A0. CP Control Block Relationships



## SVC INTERRUPTIONS

When an SVC interruption occurs, the SVC interruption routine (DMKPSASV) is entered. If the machine is in problem mode, DMKPSASV takes the following action:

- The SVC interrupt code is examined to determine if the interrupt was the result of an ADSTOP SVC code X'B3'. If it was, the message ADSTOP AT XXXXX is sent to the user's terminal, the overlaid instruction is replaced, and the virtual machine is placed in console function mode via DMKCFMBK; otherwise, the virtual machine's mode (BC or EC) is determined.
- If the virtual machine was in EC mode or its page 0 was not in real storage, then all general and floating-point registers are saved, the user's VMELOK is flagged as being in an instruction wait, and control is transferred (via GOTO) to DMKPRGRF to reflect the interrupt to the virtual machine.
- If the virtual machine was in BC mode and if his page 0 is in main storage, then an appropriate SVC old PSW is stored in his page 0 and the interrupt is reflected to the virtual machine, bypassing unnecessary register saving. If the new virtual PSW indicates the wait state, all registers are saved in the VMBLOK and control transfers to DMKDSPB for PSW validation.

If the machine is in supervisor mode, the SVC interruption code is determined and a branch is taken to the appropriate SVC interruption handler.

SVC 0: Impossible condition or fatal error. The SVCDIE routine initiates an ABEND by going to the DMKDMPDK routine.

SVC 4: Reserved for IBM use.

SVC 8: Link request (transfer control from calling routine to called routine specified by register 15). The SVCLINK routine sets up a new save area, and then saves the caller's addressability (register 12) and save-area address (register 13), and the return-address (from the SVCOPSW) in the new save area. If the called

routine (specified by register 15) is within the resident CP nucleus, SVCLINK places its address in register 12 and branches directly to the called routine. If the called routine is in a pageable module, a TRANS is performed on register 12 to ensure that the page containing the called routine is in storage. Upon return from the TRANS, the real address of the pageable routine is placed in register 12 and SVCLINK branches to the called routine. The real storage location of DMKCPPE is the end of the resident CP nucleus. Any modules loaded at a higher real storage address are defined as pageable modules.

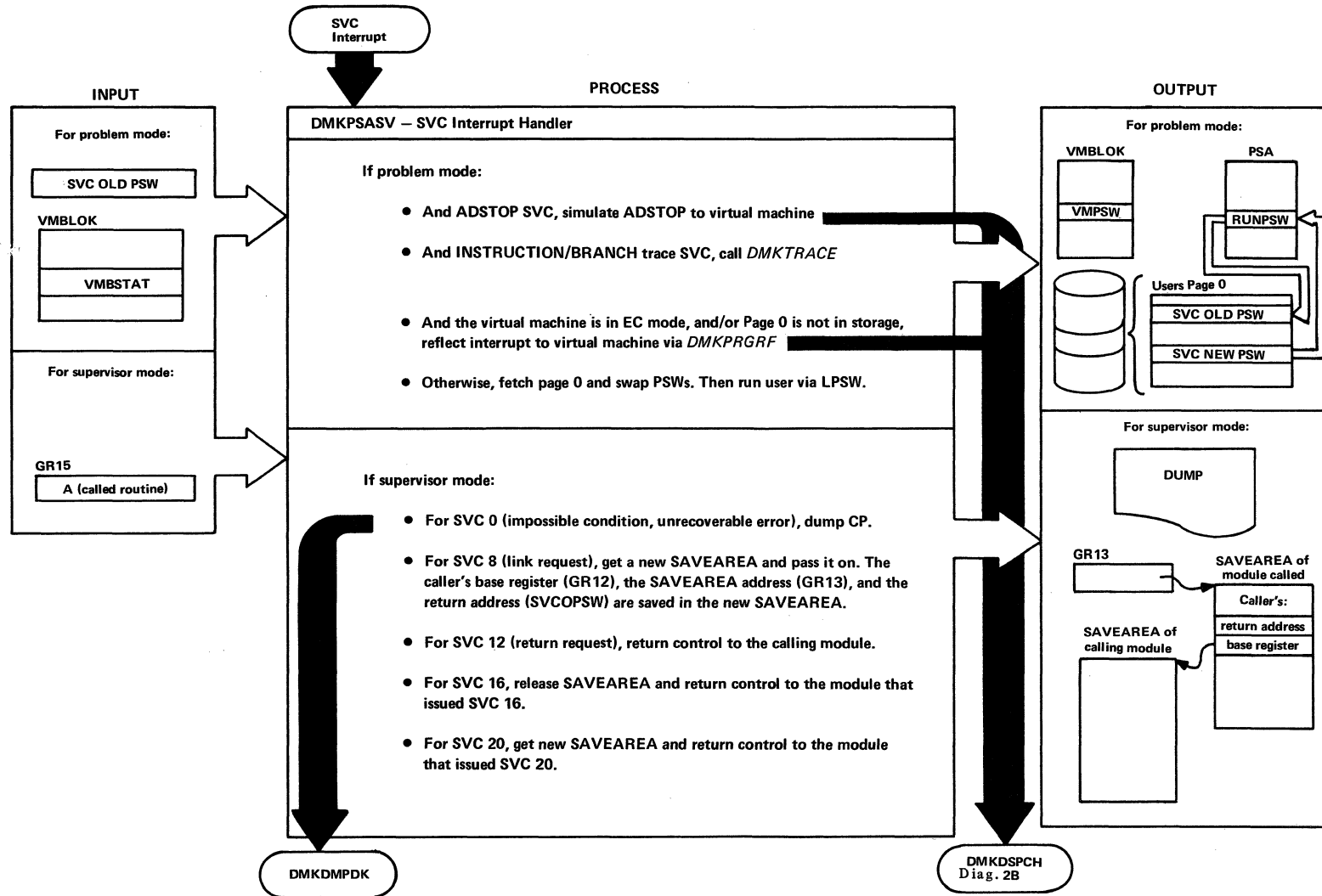
SVC 12: Return request (transfer control from called routine to calling routine). The SVCRET routine is invoked. If the routine which issued the SVC 12 is in the pageable module DMKPTRUL, then DMKPGSUL is called to unlock the page. SVCRET then restores registers 12 and 13 (addressability and save area address saved by SVCLINK), places the user's return address (also saved in the area) back into the SVCOPSW, and returns control to the calling routine by loading the SVCOPSW.

SVC 16: Release current save area from the active chain (remove linkage pointers to the calling routine). The SVCLSE routine releases the current save area by placing the address of the next higher save area in register 13 and returns control to the current routine by loading the SVCOPSW. This SVC is used by second level interrupt handlers to bypass returning to the first level handler under specific circumstances. The base address field (register 12) in the save area being released is examined to determine if the bypassed routine is in a pageable module. If so, DMKPTRUL is called to unlock the page.

SVC 20: Obtain a new save area. The SVCGET routine places the address of the next available save area in register 13 and the address of the previous save area in the save area pointer field of the current save area.

There are 35 SAVEAREAs initially set up by DMKCPINT for use by the SVC linkage handlers. If the supply of available save areas drops to zero, the linkage handlers call DMKFREE to obtain storage for additional save areas.

Diag. 1B1. SVC Interrupt Handler



## EXTERNAL INTERRUPTIONS

When an external interruption occurs, the external interruption handler (DMKPSAEX) is entered.

### TIMER INTERRUPT

If DMKPSAEX is entered because of a timer interrupt, the machine mode must be determined. If the machine was in WAIT state, control is transferred to DMKDSPCH which becomes idle until another interrupt occurs. If the machine is in problem mode, the address of the current user's VMBLOK is obtained from RUNUSER. The user's current PSW (VMPSW) is updated from the external interruption old PSW, the address of the current VMBLOK is placed in register 11, and control is transferred to DMKDSPCH. For additional information about timers see the section "Virtual Timer Maintenance".

### EXTERNAL INTERRUPT

If DMKPSAEX is entered because of the operation of the console interrupt button (INTERRUPT), the following steps are taken:

1. The current system operator's VMBLOK (DMKSYSOP) is referenced.
2. His virtual machine is disconnected.

The operator can now logon from another terminal. The operation of the console interrupt button is used to implement an alternate operator's console. For a description of the processing of the EXTERNAL command refer to module DMKCPB.

## PROGRAM INTERRUPTIONS

When a program interruption occurs, the program interruption handler (DMKPRGIN) is entered. Program

interruptions can result from:

- Normal paging requests.
- A paging request by a virtual machine in EC mode (virtual relocation).
- Privileged instructions.
- Program errors.

DMKPRGIN determines the cause of the interruption by examining the interruption code.

### NORMAL PAGING REQUESTS

If the program interrupt is caused by a normal paging request (it is not from a virtual machine that is running in EC mode with translation on), DMKPRGIN determines whether a segmentation error (a segment of the program occurred; if so, an invalid address interruption code is set, and the interruption is reflected to the user's virtual machine supervisor. If a segmentation error has not occurred, the user's current PSW is updated from the program old PSW (PROPSW), the address of the current VMBLOK is placed in register 11, and DMKPTRAN is called to obtain the required page. When the paging operation is completed, control is returned to DMKDSPCH.

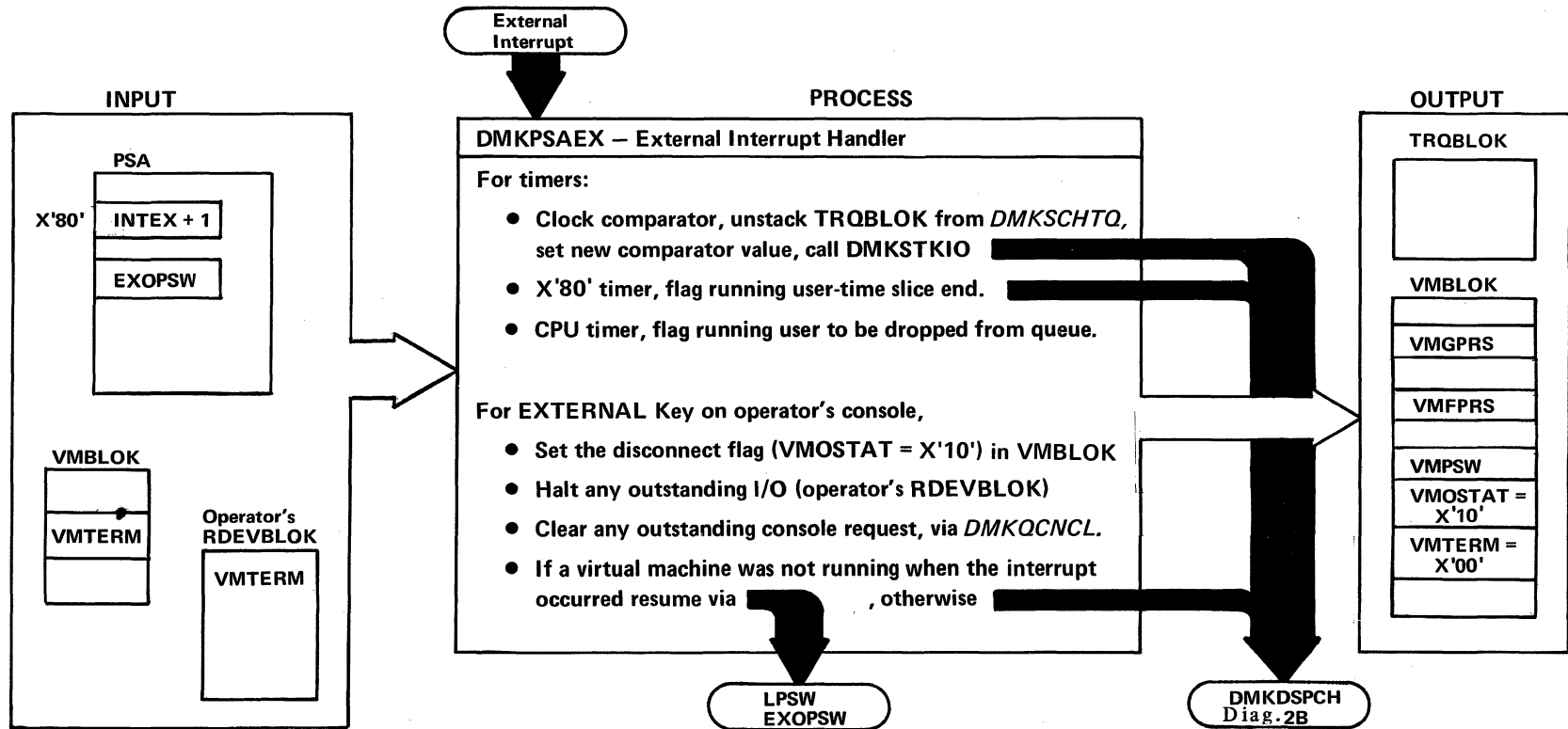
The functions of paging are divided into three categories: the management of virtual storage, the management of real storage, and the management of auxiliary storage (DASD paging devices).

### Virtual Storage Management

When operating in the relocate environment provided by CP, each user's virtual storage space is described by two sets of tables.

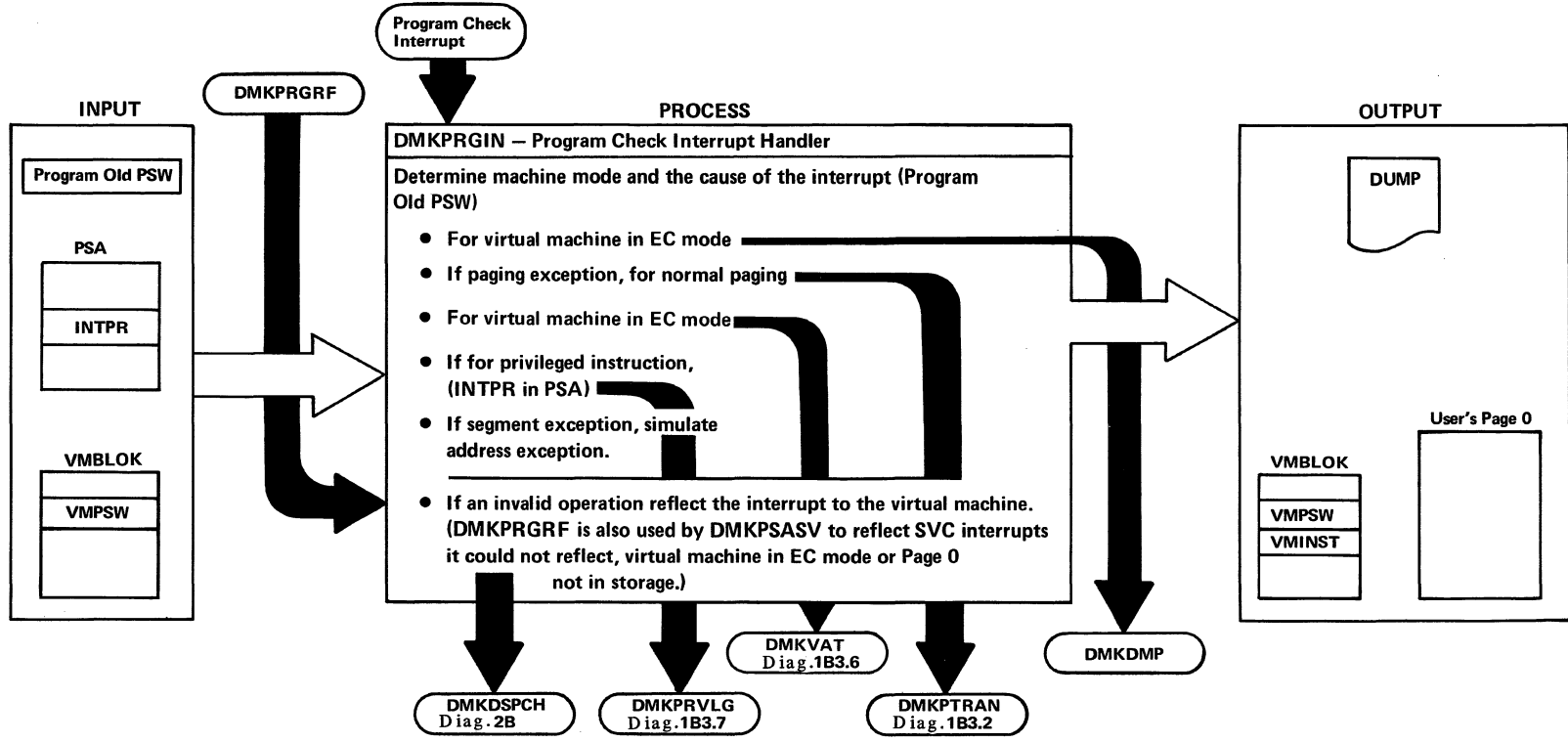
- One set, the segment and page tables, describes the location and availability of any of the user's virtual pages that may be resident in real storage. Locations in these tables are indexable by virtual address, and the entries contain index values that reference corresponding real storage addresses. In

Diag. 1B2. External Interrupt Handler

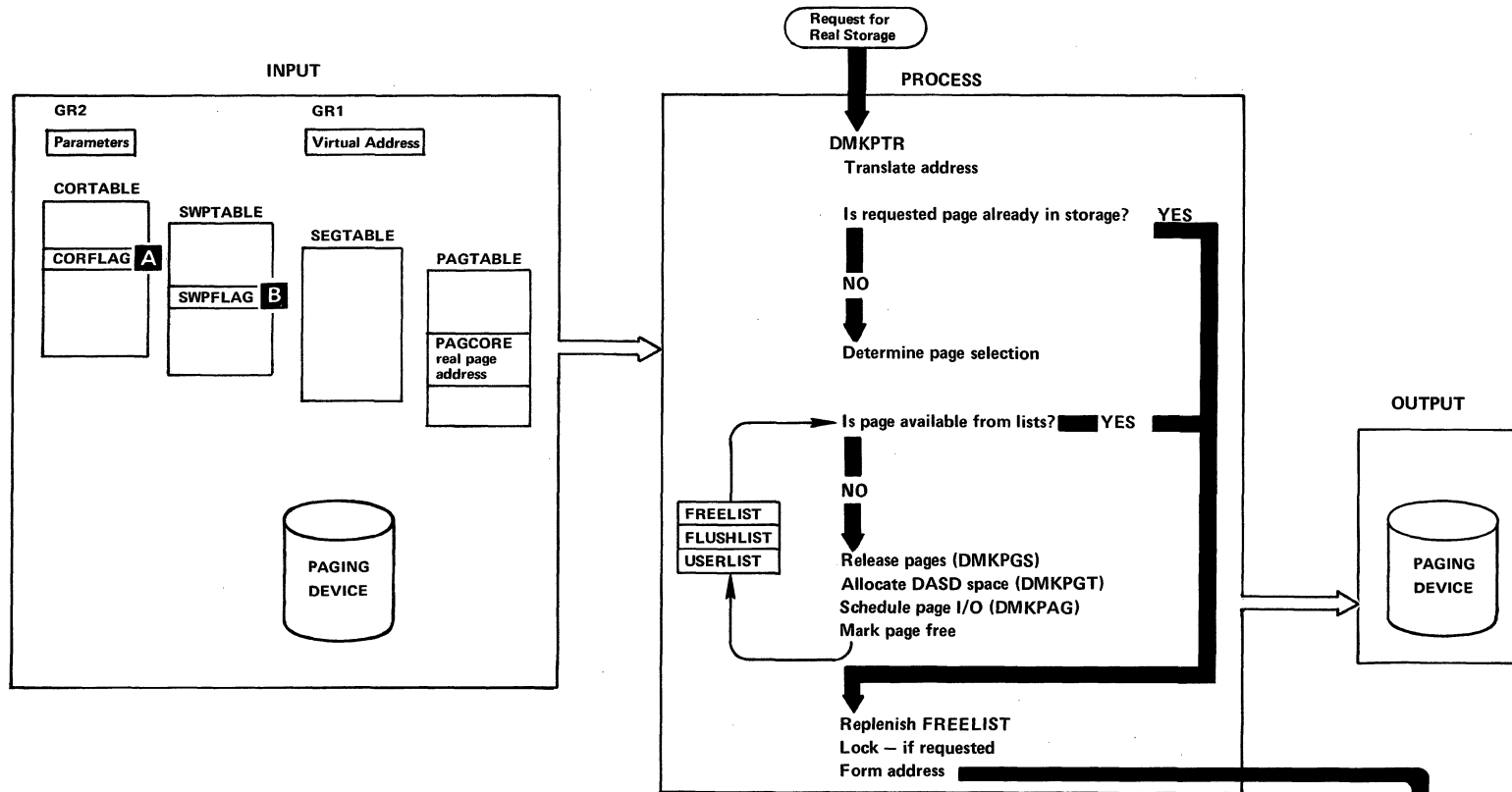




Diag. 1B3. Program Interrupt Handler

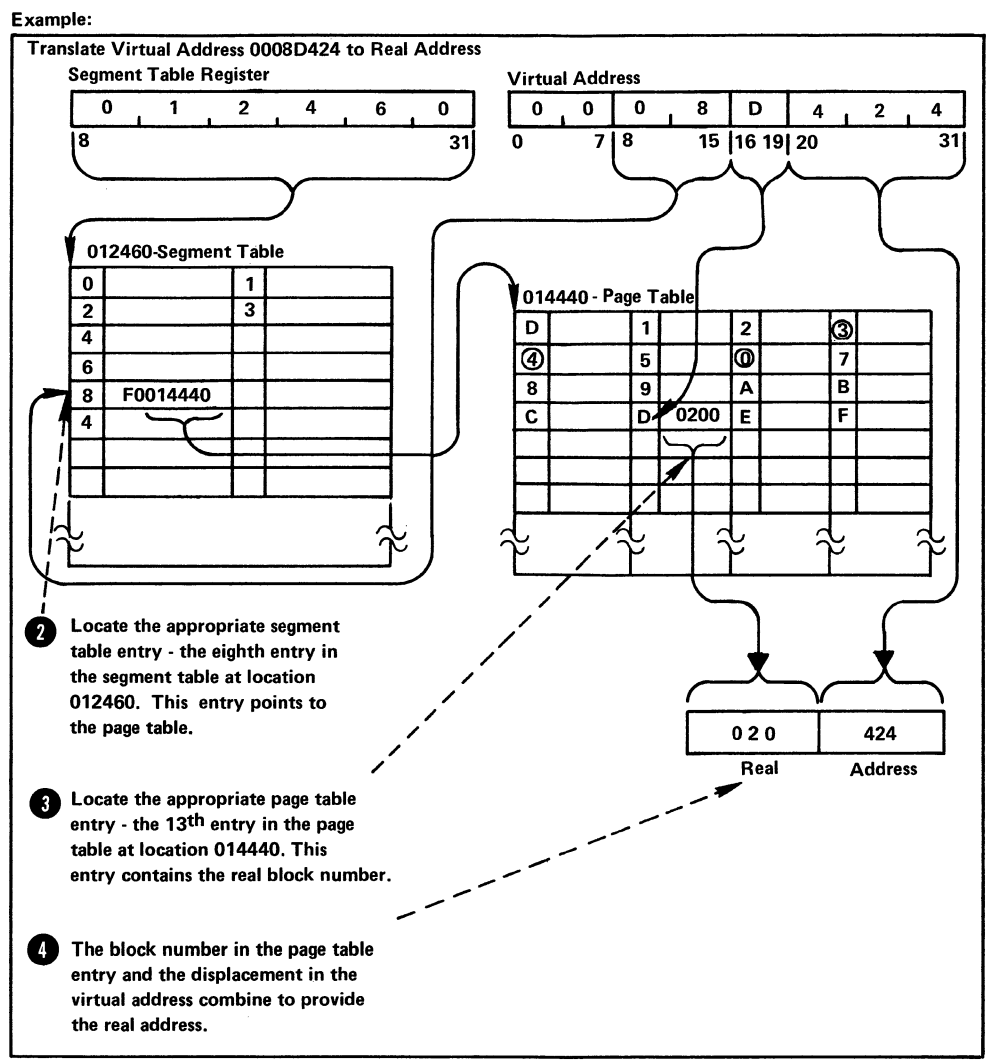
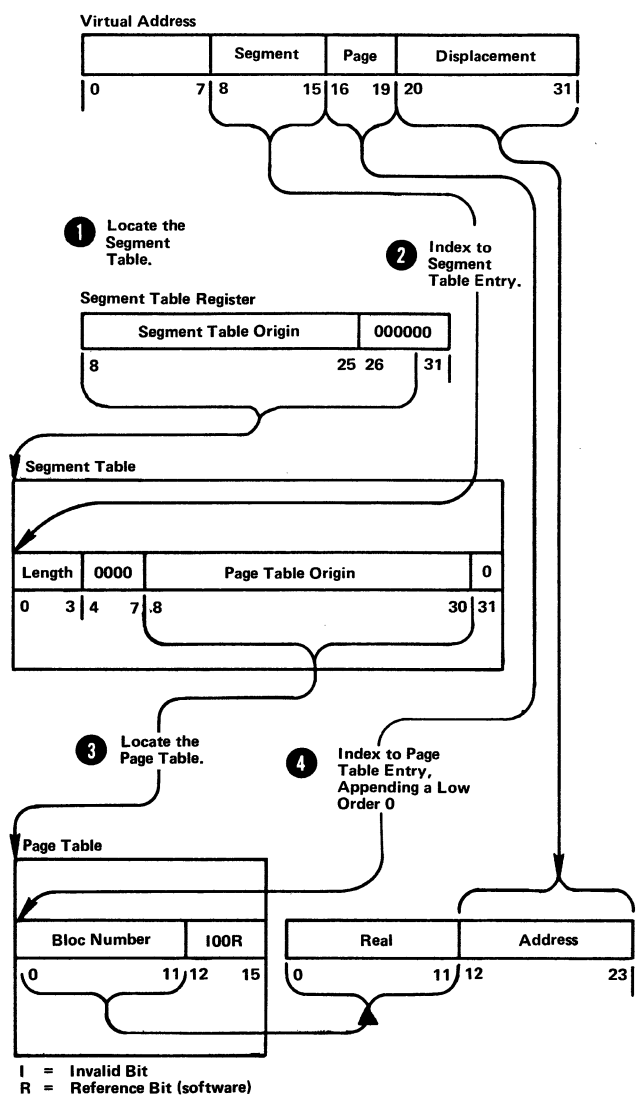


Diag. 1B3.0. Paging Overview



A Bits defined for CORFLAG			B Bits defined for SWPFLAG		
CORIOLOCK	EQU X'80'	Page locked for I/O	SWPTRANS	EQU X'80'	Page in transit
CORCFLCK	EQU X'40'	Page locked by console function	SWPRECMP	EQU X'40'	Page permanently assigned
CORFLUSH	EQU X'20'	Page is in flush list	SWPALLOC	EQU X'20'	Page enqueued for allocation
CORFREE	EQU X'10'	Page is in free list	SWPSHR	EQU X'10'	Page shared
CORSHARE	EQU X'08'	Page is shared	SWPREF1	EQU X'08'	1st half page referenced
CORRSV	EQU X'04'	Page is reserved	SWPCHG1	EQU X'04'	1st half page changed
CORDISA	EQU X'01'	Page disabled - not available	SWPREF2	EQU X'02'	2nd half page referenced
			SWPCHG2	EQU X'01'	2nd half page changed

Diag. 1B3.1. Virtual to Real Address Translation



addition, each table entry contains an indication of whether the corresponding virtual page is available to the user in real storage. These tables are referenced directly by the DAT feature when the user's program is running.

- The second set of tables is a map of the locations of the user's pages on the DASD devices that comprise the system's paging or auxiliary storage. The DASD addresses in these tables can either represent the source of a page of virtual storage (the location to which a page may be moved, if necessary) or a dummy address, indicating that the given page has not yet been referenced, and thus has a value of binary zeros.

The tables are arranged in a format indexable by virtual storage address. In addition to containing the address of a page, each entry contains flags and status bytes that indicate such information as:

- The storage protection keys to be assigned to the page when it is made resident.
- Whether the page is currently on its way into or out of the system (in transit), etc.

These tables, called swap tables, are not referenced directly by the hardware as are the page and segment routines tables, but are used by paging management to locate user pages that are needed to execute a program.

Virtual storage management is done by the technique known as demand paging. This means that a page of virtual storage is not 'paged in' from its DASD auxiliary store slot until it is needed for execution. CP does not determine the pages required by a user before he is run. A demand for a page can be made either implicitly or explicitly.

- An implicit demand is made when a user program attempts to reference a page that is not available in real main storage. This attempt causes a program interrupt with the interrupt code indicating a page or segment exception. Upon recognition of this condition, control is passed to the paging manager to obtain a page of real main storage and to bring in the desired page.

- An explicit request for virtual storage can be made by the control program (for example, in the course of translating a user's channel program). If, in the process of translation, the control program encounters a CCW that addresses a page that is not resident in real storage, a call is made to the paging manager to make the referenced page resident.

While the requested page is being fetched, the requesting user is unable to run; however, it may be possible to run other tasks in the system, and the control program runs these while that page is being paged in. When the requested page is resident, the user can be run and is dispatched in his turn.

In addition to obtaining pages by demand, users implicitly or explicitly release pages of their virtual storage space. Part of the space may be explicitly released from both real and virtual storage via a diagnose instruction which indicates to the control program those pages that are to be released. An entire virtual storage is released when a user IPLs a new operating system or logs out from the system.

The VM/370 control program itself also has virtual storage associated with it. This space is used to contain the control program (some parts of which need not always be resident in real storage), and is also used for virtual storage buffers for spooling and system directory operations. Although the control program makes use of virtual storage space for its execution, it does not run in relocate mode itself. Thus, nonresident modules must be completely relocatable.

#### Real Storage Management

It is the function of real storage management to efficiently allocate the system's page frames of real storage to satisfy the demands for virtual pages made by the system's users. Efficiency of allocation involves a trade-off; the paging manager utilizes only enough CPU time allocating to ensure that:

1. The set of virtual storage pages which are resident represent those pages that are most likely to be used.

2. The number of cycles available to execute user programs is sufficient.

Inefficiency in the first area causes a condition known as thrashing, which means that highly used pages are not allowed to remain resident long enough for useful work to be performed by or on them. Thrashing could be aggravated by the paging manager's page selection algorithm or by a dispatcher that attempts to run more tasks than the system can handle (the sum of their storage requirements exceeds the real paging space available in the system). Thus, the paging manager must keep statistics on system and user paging activity and make these statistics available to the dispatcher so that a potential thrashing condition may be detected and prevented.

Inefficiency in the second area causes an unacceptable ratio of control program overhead to user program time, and in extreme case may cause the control program to utilize excessive CPU time. In order to understand how allocation is determined by the VM/370 control program, it is first necessary to describe the way in which the inventory of real storage page frames is described to the system.

Each page frame (4096 byte blocks) of real storage in the system is in one of two basic states: not-pageable or pageable. A not-pageable page must remain resident in real storage for some finite period of time; thus, the page frame cannot be taken from its current owner in order to give it to someone else. Pages can be not-pageable either permanently or temporarily, depending on their use.

- Temporary locks usually occur when an I/O operation has been initiated that is moving data either to or from the page, and the page must be kept in real storage until the operation has completed.
- A page can also be temporarily not-pageable if it contains a nonresident control program routine that is active.

In addition, a page can be not-pageable through use of the LOCK command. Pages locked in this fashion are permanently resident until they are explicitly unlocked by the UNLOCK COMMAND. Pages that are usually considered permanently not-pageable are those that

contain the resident portion of the control program and those that contain the system's free storage area in which control blocks, I/O buffers, etc. are built.

CORTABLE: The data area that is used by the page management routines to control and allocate real storage is the CORTABLE. Each page frame of real storage has a corresponding entry in the CORTABLE, and since the table entries are fixed length and contiguous, the entry for any given real page frame may be located directly by indexing into the table. Each entry contains pointers that indicate both the status and ownership of the real page which it represents. Some pointers are used to link page table and swap table entries to the real page (and thus establish ownership), while others are used to link the entry into one of several lists that the paging routines use to indicate the page's status and availability for paging. A given CORTABLE entry may appear on one of three lists if its real page is available for paging; however, if the page is locked or in transit, its entry is not in any list and is not referenced when available page frames are being searched for swap candidates. The lists are known as the FREELIST, the FLUSHLST, and the USERLIST, and they represent various levels of page availability.

- The FREELIST contains page frames that are immediately available for assignment to a requesting user. The virtual storage pages for which they were last used have either been released by their owners or they have been paged out to auxiliary storage. Requests for real storage are always satisfied from the FREELIST. If the list has been depleted, the requestor waits until a new page frame becomes available as the result of a virtual storage release or a swap-out.
- The FLUSHLST contains page frames that belong to those users that have been dropped from an active DISPATCHING queue. The FLUSHLST is the first place that the page frame selection routine looks to find a page to swap out or to assign to the FREELIST for a user who requires real storage space.
- The USERLIST contains the CORETABLE entries for all other pageable pages in the system that belong to active users.

Requests For Real Storage Pages

Requests for real storage fall into two general categories; those that are requesting space for a page of virtual storage, and those (such as requests for CP work space) that need the real page for their own use. The former, more general case is discussed first, since the latter case is a subset of the first.

The main page manager routine, DMKPTRAN, maps a request for a specific user's virtual storage address into a page of real storage. This requires that:

- The virtual page be read in.
- The necessary tables be updated to show the proper status of the page.

DMKPTRAN requires that the caller supply only the virtual address to be translated and any options that apply to the page to be located. Most calls are made via the TRANS macro, which sets up the necessary parameters, determines if the required page is resident, and calls DMKPTRAN if it is not.

When DMKPTRAN receives control, it first tests to see if the requested page is resident. This is done via the LRA hardware translation feature. If the page is resident, the routine locks the page if requested and exits to the caller. If the LRA indicates that the page is unavailable, it is still possible that the required page is resident. This occurs if the page has been placed on the FREELIST but has not been assigned to another user. When the page swap routine removes a page from a user, the unavailable bit is set in the corresponding page table entry; however, the real main storage index for the page is left unchanged. The page table entry is set to zero only when the corresponding page is actually assigned to another user. Thus, if DMKPTRAN finds the page unavailable, a further test is made on the page table entry to see if the page can be reclaimed. If the entry is not zero (aside from the unavailable bit), the CRTABLE entry for the page is removed from the FREELIST and the page is returned to the calling user.

If the page table entry corresponding to the virtual page requested is zero, the required page is not in real storage and must be paged in. However, it is

possible that the page is already on its way into main storage. This condition is indicated by a flag in the SWPTABLE entry for the virtual page. The DMKPAGIO routine maintains a queue of CPEXBLOCKS to be dispatched when the pending page I/O is complete. The CPEXBLOCK for the page in transit is located and a new CPEXBLOCK, representing the current request, is chained to it.

Before exiting to wait for the paging operation to complete, DMKPTRAN checks to see if the deferred return (DEFER option) has been specified. If it has not, DMKPTRAN returns to the caller. If the DEFER option has been requested, DMKPTRAN exits to the dispatcher to wait for page I/O completion. When the requested page has been read into real storage, the list of CPEXBLOCKS are unstacked FIFO to satisfy all requests for the page that arrived while it was in transit.

If a page is not in transit, a page frame of real storage must be allocated to fill the request. Before the allocation routine is called, a test is made to see if the caller wishes the return to his routine or to be delayed until after the requested page is available. If the DEFER option is not requested, DMKPTRAN returns to the caller after first building and stacking a CPEXBLOCK that allows processing of the page request to be continued the next time the dispatcher (DMKDSPCH) is entered.

DMKPTRAN next calls the FREELIST manager (DMKPTRFR) to obtain the address of the next available CORTABLE entry. DMKPTRFR maintains a FIFO list of the CORTABLE entries for those page frames that are immediately available for assignment. As DMKPTRFR releases these page frames, a check is made to see if the number of entries on the FREELIST has fallen below a dynamically maintained minimum value. If it has, the page selection routine (SELECT) is called to find a suitable page for placement in the FREELIST. The number maintained as the FREELIST threshold has a value equal to the number of users in queue1 plus the number of users in queue2 plus 1.

The FREELIST is replenished directly by users releasing virtual storage space. The page-out routine DMKPGSPD calls DMKPTRFT to place released pages directly on the FREELIST. However, most replenishment is done via the page selection routine, SELECT. SELECT is called by DMKPTRFR when the FREELIST count falls below the current minimum, or when a user page is reclaimed from

the FREELIST. In either case, the selection algorithm attempts to find a page to swap to auxiliary storage. The highest priority candidates for a swap are those pages whose CORTABLE entries appear on the FLUSHLST. SELECT attempts to take a flushed page before it takes a page from an active user. If such a page is found, it is checked to see if it has been changed since page-in. If not, it is placed in the FREELIST by DMKPTRFT; otherwise, it is scheduled for a swap-out by dequeuing the CORTABLE entry from the FLUSHLST, constructing a CPEXBLOK for dispatching after I/O completion, and exiting to DMKPAGIO via a GOTO. After the paging I/O is complete, the entry is placed on the FREELIST via a call to DMKPTRFT.

If the FLUSHLST is exhausted, SELECT must take a page from an active user by examining the pages represented by the entries in the USERLIST to locate the least recently used user page. This list is scanned from top to bottom, and each page is tested to see if its hardware referenced bits have been set. If a page has been referenced, its bits are reset and it is queued to the end of the USERLIST. This process is continued until either an unreferenced page is found or the list is exhausted. An unreferenced page is immediately selected. However, if the list is exhausted, it is rescanned from the top. An unreferenced page is always found; in the worst case it is the first one tested on the USERLIST at initial entry. However, if this occurs, it indicates that the rate of entry to SELECT is too low to permit differentiation between high and low usage pages.

Once a page has been selected and its page-out is scheduled, control is returned to DMKPTRFR, which then passes control back to DMKPTRAN with the address of the CORTABLE entry that was allocated. In most cases, page-outs are completely overlapped with page-ins. Approximately one half of all page-ins require a corresponding page-out.

Once a real page has been assigned, DMKPTRAN checks to see if a page-in is required. It usually is, and the DASD address of the virtual storage page must be obtained from the user's swap table entry and the I/O operation scheduled. However, if the page has not yet been referenced (as indicated by a DASD address of zero), the real main storage page is set to zero. After the page-in operation has been queued, DMKPTRAN exits to the paging I/O scheduler (DMKPAGIO) which initiates

the paging operation and exits to the dispatcher (DMKDSPCH) to await the interrupt.

After the required page has been read in or set to zero, DMKPTRAN queues the appropriate CORTABLE entry to the end of the USERLIST, where it eventually is available for page selection. After developing the real storage address that corresponds to the requested virtual address, DMKPTRAN tests to see if the caller has requested that the page be locked. If LOCK is requested, the CORTABLE entry is de-queued from the USERLIST and is not available for selection. A resident page can also be locked by removing it from the USERLIST. In addition, a LOCK count is maintained in the CORTABLE entry so that when all locks have been satisfied the page can again be made available for paging (see PAGUNLOK).

Some requests for main storage pages are handled differently than the general case of virtual-to-real storage mapping. In particular, it may be necessary for CP to obtain additional free storage for control blocks, I/O lists, buffers, etc. This is handled by the free storage manager, which makes a direct call to DMKPTRFR to obtain the needed storage. Usually this storage is immediately available (due to the page buffering technique previously described). However, if the FREELIST is exhausted, the request for free storage is recognized as a high priority call and queued first on the list of those waiting for free pages.

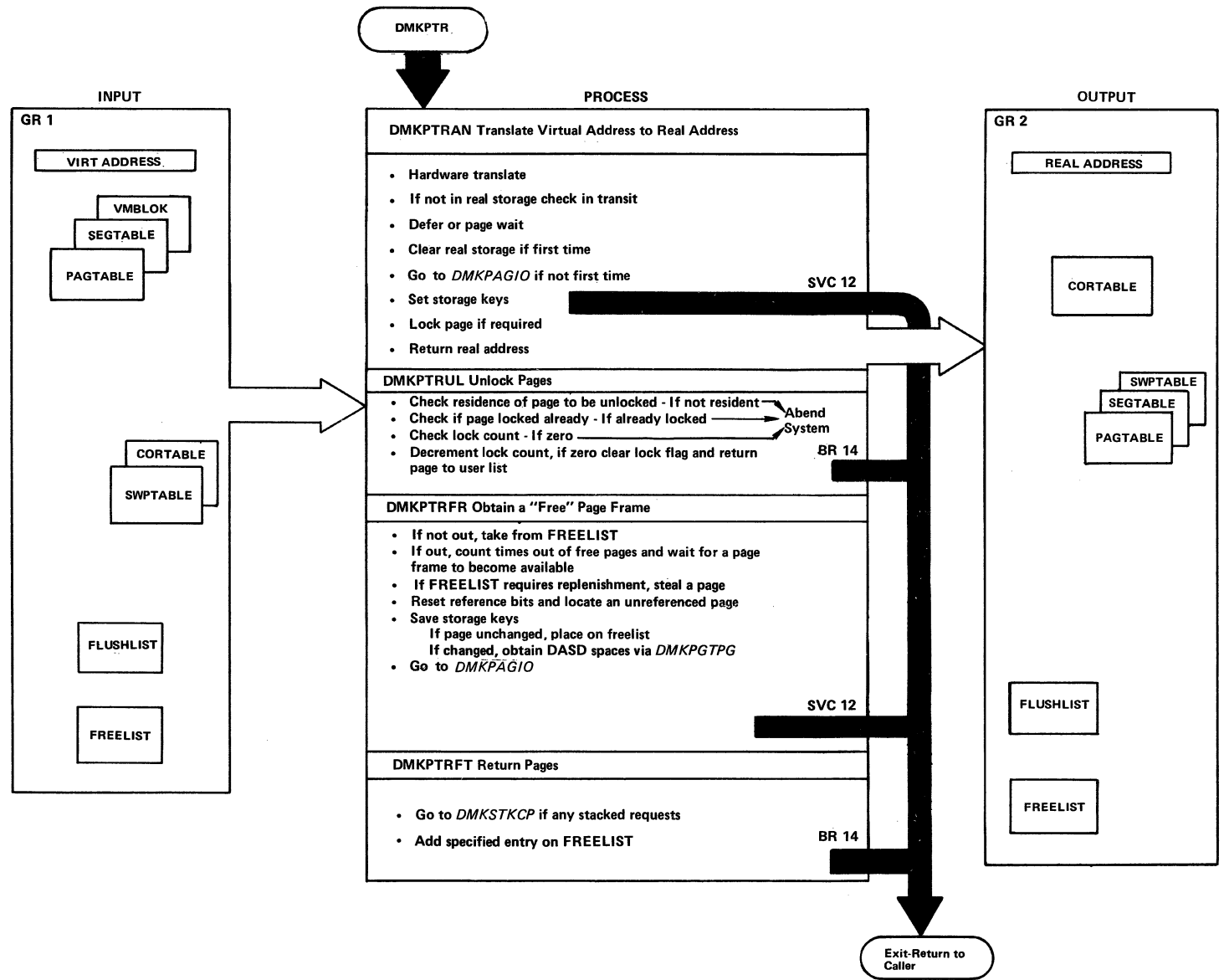
The real storage manager (DMKPTR) accumulates paging statistics which are used by the scheduler (DMKSCH) to project user storage requirements. A count of page-reads and page-writes is kept in each user's VMBLOK; the corresponding total counts for the system are kept in DMKPSA. A running total of the number of pages a user has resident, at each instance of page-read, is kept in the VMBLOK. A count of the number of times a user enters page-wait, because a page has been stolen from him, is also kept in the VMBLOK. The section entitled "Controlling the Depth of Multiprogramming" under the heading "Dispatcher/Scheduler" describes the use to which the scheduler puts these counts.

VM/370 Virtual=Real Function: The VM/370 Virtual=Real function involves the mapping in a one-for-one correspondence of a virtual machine storage area with an equivalent real storage area. For instance, virtual

page 1 is in real page frame 1 and virtual page 20 is in real page frame 20. Virtual page 0, since it cannot occupy real page 0, is relocated to be at the end of the virtual storage space.



Diag. 1B3.2 Paging, Provide Real Storage Area





The VM/370 control program nucleus is altered at system generation to support the Virtual=Real function. Users with Virtual=Real (specially identified in the directory) can then log on and use the space reserved for this function. That space can be used by only one virtual machine at a time. Two virtual machines with the Virtual=Real capability cannot occupy the same space at the same time.

The Virtual=Real function is primarily used so that the virtual machine may bypass the control program's CCW translation. This is possible because I/O from a virtual machine occupying a Virtual=Real space contains a list of CCWs whose data addresses reflect the real storage addresses. The restriction in this situation is that the virtual machine does not perform I/O into page 0 since this would perform a data transfer into real page 0. At the same time, it is assumed, and cannot be checked, that the virtual machine will also not attempt to do I/O beyond the bounds of its virtual addressing space. To do so would cause the destruction of either the VM/370 control program nucleus, which resides beyond the virtual machine space, or another user's page.

The bypassing of CCW translation for the virtual machine occupying the Virtual=Real space is only invoked after the virtual machine has executed the SET NOTRANS ON function. This function can only be issued by the virtual machine occupying the Virtual=Real space. The function initiates the bypass of CCW translation. This function is automatically turned off if the virtual machine performs an explicit reset, or an implied reset by performing a virtual IPL. During virtual machine IPL, it is required that I/O be performed into page 0. For this reason, normal virtual IPL simulation assumes CCW translation in effect in order to accomplish the full simulation. Once the IPL sequence has completed, the CCW translation function can be bypassed by issuing the SET NOTRANS ON command.

When the virtual machine demands a page through normal use of the control program's page tables, the paging routine recognizes the Virtual=Real capability. It then assigns the virtual page to the equivalent real page frame and does not perform a paging operation, since all these pages are resident and are never swapped out.

Note: The virtual machine running with Virtual=Real is still run in System/370 relocate mode.

Virtual 270X lines and sense operations from the virtual machine do not use the Virtual=Real feature. These invoke CCW translation for the virtual enable/disable lines and the transfer of the sense bytes.

The UNLOCK command has an operand called VIRT=REAL and essentially releases the Virtual=Real area for normal system paging use. Once the area has been released, it can only be reclaimed by an IPL of the VM/370 System. The size of the Virtual=Real area is an installation specification that is part of the special nucleus generation procedure that is outlined in the VM/370 Planning and System Generation Guide. The size of the area must be large enough to contain the entire addressing space of whatever virtual machine wishes to occupy that space. A virtual machine can use a smaller space than is provided but cannot use a larger space without regenerating the VM/370 control program nucleus.

#### DASD Storage Management

Any user virtual storage pages that have been referenced but are not resident in real storage must be kept on the DASD paging device. DASD page space is assigned only when the page is selected for a page-out. Certain DASD pages may also be marked read-only. Thus, the DASD address slot initially associated with the page should be considered to be the source of the page only. If the page is changed after it has been read into real storage, a new slot must be obtained when it is paged out. Examples of read-only pages are those which contain portions of pageable saved systems and pages which are part of a system SPOOL file. Slots can be reassigned when DMKPTRAN finds that it must swap a page out to a movable head DASD device. In this case, the old slot is released and the new slot is obtained.

SLOT ALLOCATION: If a new slot is required, the DMKPGT is called to supply the address of an available slot. DMKPGT maintains a chain of cylinder allocation maps for each cylinder that has been assigned for either virtual storage or spool file paging. The allocation

chains for spooling are kept separately from those used for paging so that they can be checkpointed in case of a system failure. However, in other respects they are the same. The allocation blocks for a given volume are chained from the RDEVBLK for the device on which the volume is mounted. The chains of cylinder and slot allocation blocks are initialized by DMKCP1. Each block on an allocation chain represents one cylinder of space assigned to paging, and contains a bit map indicating which slots have been allocated and which are available. Each block also has a pointer to the next allocation block on the chain, a cylinder number, and a record count. DMKPGT searches this list sequentially until an available slot is found; its DASD address is then determined and passed back to the calling routine. If DMKPGT cannot find a cylinder with a de-allocated slot, it enters the cylinder allocation phase described next. When an available cylinder is found, it constructs a page allocation block for this cylinder and allocates a page to the caller.

**CYLINDER ALLOCATION:** DMKPGT controls the paging and spooling I/O load of the system by allocating cylinders evenly across all available channels and devices. In order for a device to be considered available for the allocation of paging and spooling space:

- Its volume serial number must appear in the system's owned list.
- It must have at least one cylinder of temporary space marked as available in the cylinder allocation block which is located on cylinder 0, head 0, record 3.

At system initialization time, CPINIT reads in the allocation records for each volume and constructs the chains of device allocation blocks from which DMKPGT allocates the cylinders. In managing the cylinder allocation, DMKPGT takes three factors into consideration: device type, device address, and possible status as a preferred paging device.

A request for a cylinder of virtual storage page space is satisfied by allocating on a preferred paging device, provided that one exists on the system and that it has page space available. Preferred paging devices are specified by the installation at system generation time, and generally should be devices on which excessive seek times does not occur. A typical

preferred paging device would be the IBM 2305 Fixed Head Storage Facility. If the 2305 is assigned as a preferred device, it is possible to allocate some of its space for other high priority data files without excessively degrading paging. An example of such usage would be for high activity read-only saved system pages that are not shared in real storage, and high activity system residence disks.

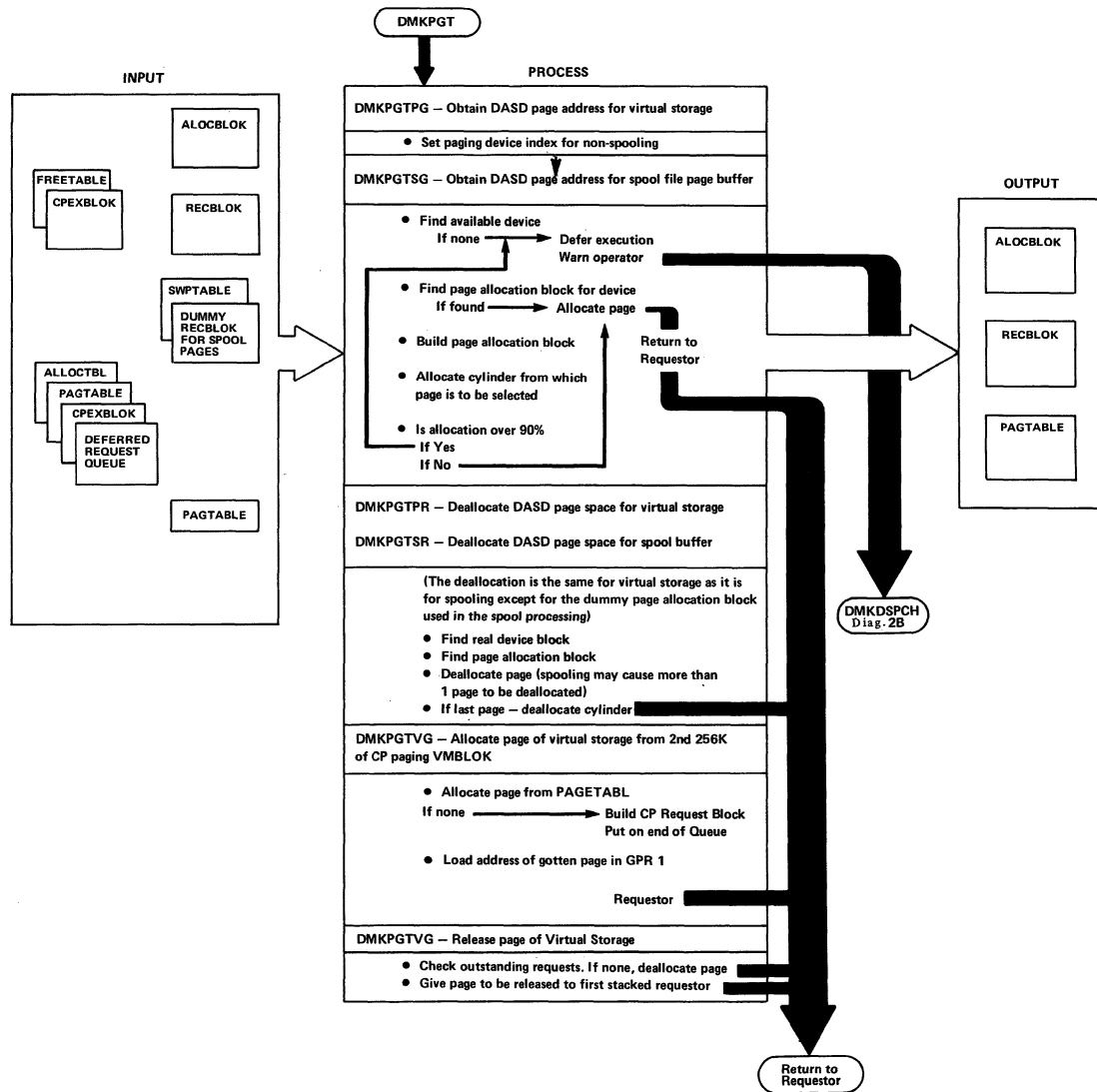
It is also possible to designate moveable head DASD devices such as the 3330 and 2314/2319 Direct Access Storage Facilities as preferred paging devices. The module(s) so designated should not be required to seek outside of a relatively narrow cylinder band around the center of the paging areas. It is advisable to share the access arm of a moveable head preferred paging device with only the lowest usage data files.

If one or more preferred devices are defined on the system, CP allocates all of the page space available on these before it allocates on any other available owned volumes. Within the class of preferred devices, space is allocated first on the fastest devices, and among these on a round robin basis across channels and devices. Allocation on nonpreferred devices is spread out in the same manner. Cylinders for spooling space are not allocated from preferred devices. Allocation on a given device is done from the relative center of the volume outward, a cylinder at a time in a zig-zag fashion in an attempt to minimize seek times.

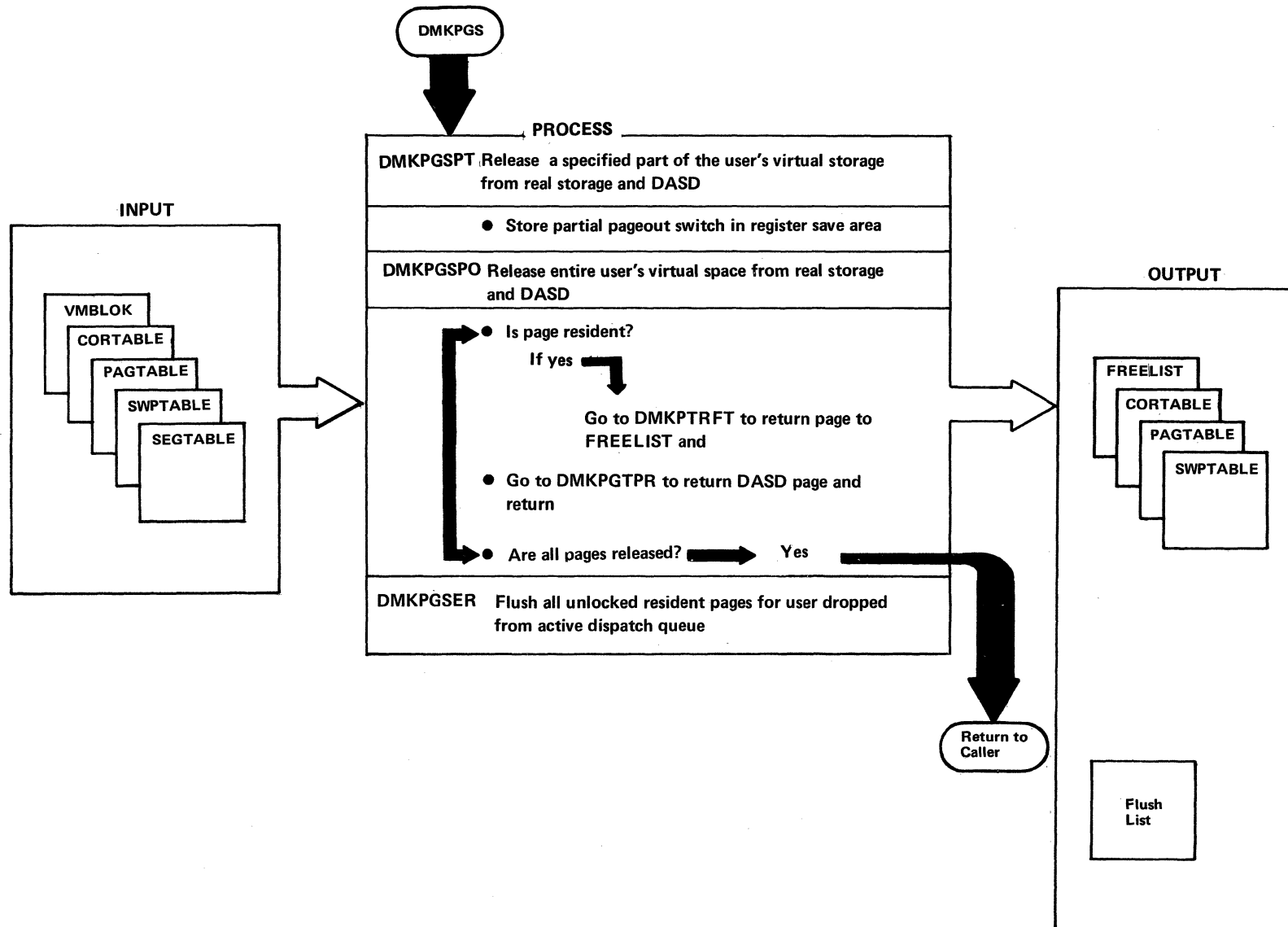
When a request to allocate a slot for virtual storage paging is received by DMKPGTGT and the slot must be allocated on a moveable-head (2314/2319 or 3330) device, a cylinder and slot is selected in the following manner:

1. An attempt is first made to allocate a slot on the cylinder at which the arm on the selected device is currently positioned.
2. If slots are not available on the current cylinder, an attempt is made to allocate on a cylinder for which paging I/O has been queued.
3. If the above conditions cannot be met, allocation is done as close to the center of the volume as is possible.

Diag. 1B3.3. Paging, Allocate DASD Space



Diag. 1B3.4. Release Virtual Machine Pages



Before DMKIOSQR is called, the queue of IOBLOKS currently scheduled on the device is examined. If paging I/O has already been scheduled on a device, the paging channel programs are slot sorted and chained together with TICs.

### Paging I/O

All input/output requests for virtual storage and spooling pages are handled by DMKPAGIO. DMKPAGIO constructs the necessary task blocks and channel programs, expands the compressed slot addresses, and maintains a queue of CPEXBLOKS for pages to be moved. Once the I/O scheduled by DMKPAGIO completes, it unchains the CPEXBLOKS that have been queued and calls DMKSTKCP to stack them for execution. DMKPAGIO is entered via a GOTO from:

- DMKPTRAN to read and write virtual storage pages
- DMKRPA to read and write virtual storage spool buffers

In any case, all that need be passed to DMKPAGIO is the address of the CORTABLE entry for the page that is to be moved, the address of a SWPTABLE entry for the slot, a read or write operation code, and the address of a CPEXBLOK that is to be stacked for dispatching after the I/O associated with the page has completed. DMKPAGIO obtains an IOBLOK and builds a channel program to do the necessary I/O, and uses the device code that is part of the page address to index into the system's OWNDLIST and locate the real device to which the I/O request should be directed. If the device is capable of rotational position sensing, the required sector is computed and a Set Sector command is inserted into the channel program. The real SIO supervisor DMKIOSQR is then called to schedule the operation on the proper device.

When the interrupt for the paging operation is processed by the primary I/O interrupt handler, the

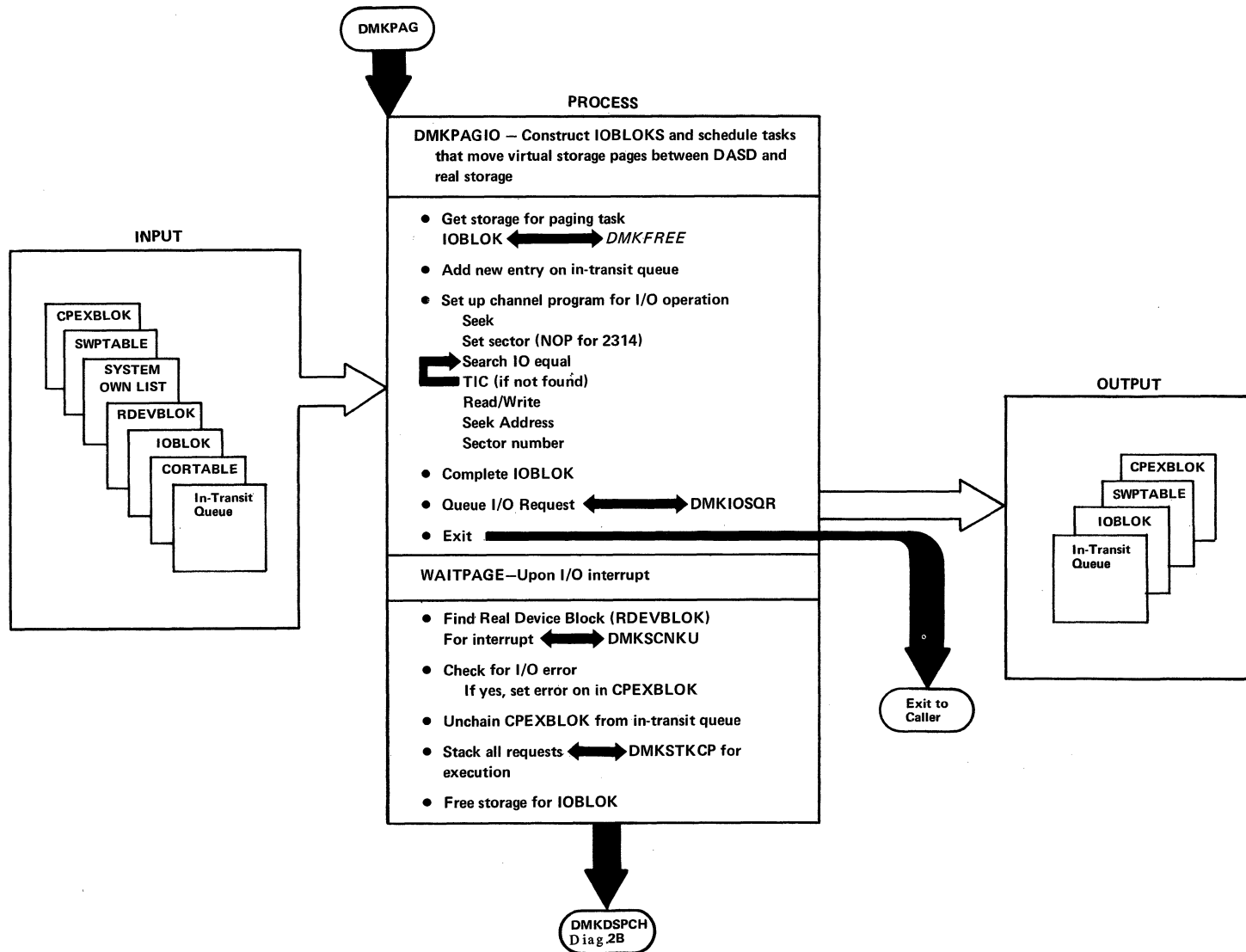
IOBLOK that controls the operation is unstacked to the interrupt return address, WAITPAGE, in DMKPAGIO. WAITPAGE then unchains the CPEXBLOKS that are queued to DMKPAGQ, and then stacks the queued CPEXBLOKS, via calls to DMKSTKCP, in the order in which they were received. The address of the real page is filled in to the appropriate page table entry and the pointers denoting the ownership of the real page are filled into the CORTABLE entry by the processing routines in DMKPTRAN. If a fatal I/O error occurred for the page, the CPEXBLOKS associated with it are flagged, and the dispatcher DMKSDPCH sets a nonzero condition code when it activates the pending task. The error recovery followed depends on the operation being performed. Paging I/O errors associated with spooling operations are discussed in the sections on "Virtual and Real Spooling", while errors associated with virtual storage paging operations are discussed later in section "Virtual Storage Paging Error Recovery".

DMKPAGIO maintains its own subpool of preformatted paging IOBLOKS. As I/O operations complete, their IOBLOKS are added to a list of available blocks; as new blocks are needed, they are taken from this list. If the list is empty, DMKFREE is called to obtain storage for a new block. DMKPAGIO also periodically calculates system paging overhead. After 200 pages have been moved (read or written), the elapsed time for the 200 pages is computed, and the paging rate is calculated in pages per second. The recent paging load, expressed as the percentage of time that more than one half of the system's pages were idle due to page-wait, is averaged with the previous load and re-projected as the expected load for the next interval.

### Virtual Storage Paging Error Recovery

Errors encountered during virtual storage (as opposed to spooling) paging operation can generally be classified as either soft or hard errors. Soft errors allow the system to continue operation without delay or degradation. Hard errors can cause noticeable effects such as the abnormal termination of user tasks (ABEND) and response degradation. Errors that are successfully

Diag. 1B3.5. Page-in, Page-out





retrieved or corrected are known only to the I/O supervisor and the I/C error retry and recording routines; they appear to the second level interrupt handlers (such as WAITPAGE) as if the original operation completed normally.

SOFT ERROR RECOVERY: An I/O error which occurs on a page swap-out is considered to be a soft error. DMKPTRAN calls DMKPGTPG to assign a different DASD page slot and the page is re-queued for output. The slot which caused the error is not de-allocated, and thus is not assigned to another user. All other uncorrectable paging errors are considered hard in that they may more drastically affect system performance.

HARD ERROR RECOVERY: Hard paging errors occur on either I/O errors for page reads or upon the condition of exhausting the system's spooling and paging space. Recovery attempted on hard errors depends upon the nature of the task for which the read was being done. If the operation was an attempt to place a page of a user's virtual storage into real storage, the operation of that particular virtual machine is terminated by setting the page frame in error to zero and placing the virtual machine in console function mode. The user and operator are informed of the condition, and the page frame causing the error is not de-allocated, thereby insuring that it will not be allocated to another user.

The control program functions which call DMKPTRAN (such as spooling, pageable control program calls, and system directory management) have the option of requesting that unrecoverable errors be returned to the caller. In this case, the CP task may attempt some recovery to keep the entire system from terminating (ABEND). In general, every attempt is made to at least allow the operator to bring the system to orderly shut-down if continued operation is impossible.

Proper installation planning should make the occurrence of a space exhaustion error an exception. An unusually heavy user load and a backed-up spooling file could cause this to happen. The operator is warned when 90% of the temporary (paging/spooling) space in the system is exhausted. He should take immediate steps to alleviate the shortage. Possible remedies that exist include preventing more users from logging on and requesting users to stop output spooling operations. More drastic measures might include the purging of low priority pool files. If the system's paging space is

completely exhausted, the operation of virtual machines progressively slows as more and more users have paging requests that cannot be satisfied and operator intervention is required.

#### VIRTUAL RELOCATION

CP provides the virtual machine the capability of using the Dynamic Address Translation of the real System/370. Programming simulation and hardware features are combined to allow usage of all of the available features in the real hardware, (that is, 2K or 4K pages, 64K or 1M segments).

For clarification, some term definitions follow:

First-level storage: The physical storage of the real CPU, in which CP resides.

Second-level storage: The virtual storage available to any virtual machine, maintained by CP.

Third-level storage: The virtual storage space defined by the system operating in second-level storage, under control of page and segment tables which reside in second-level storage.

Page and segment tables: Logical mapping between first-level and second-level storage.

Virtual page and segment tables: Logical mapping between second-level and third-level storage.

Shadow page and segment tables: Logical mapping between first-level storage and third-level storage.

A standard, non-relocating virtual machine in CP is provided with a single control register, control register zero that can be used for:

- Extended masking of external interrupts.
- Special interrupt traps for SSM.
- Enabling of virtual block multiplexing.

A virtual machine that is allowed to use the extended control feature of System/370 is provided with a full complement of 16 control registers, allowing virtual monitor calls, PER, extended channel masking, and dynamic address translation.

An extension to the normal virtual-machine VMBLOK is built at the time that an extended control virtual machine logs onto CP. This ECBLOK contains the 16 virtual control registers, 2 shadow control registers, and several words of information for maintenance of the shadow tables, virtual CPU timer, virtual TOD clock comparator, and virtual PER event data. The majority of the processing for virtual address translation is performed by the module DMKVAT, with additional routines in DMKPRG, DMKPRV, DMKDSP, DMKCDB, DMKLOG, DMKUSO, and DMKPTR. The simulation of the relocation-control instructions (that is, LCTL, STCTL, PTLB, RRB, and LARA) is performed by DMKPRV. These instructions, with the exception of LCTL and STCTL, are not available to virtual machines which are not allowed the extended-control mode.

When an extended control virtual machine is first active, it has only the real page and segment tables provided for it by CP and operates entirely in second-level storage. DMKPRV examines each PSW loaded via LPSW to determine when the virtual machine enters or leaves extended control or translate mode, setting the appropriate flag bits in the VMBLOK. Flag bits are also set whenever the virtual machine modifies control registers 0 or 1, the registers that control the dynamic address translation feature. DMKDSP also examines PSWs that are loaded as the result of interrupts to determine any changes in the virtual machine's operating mode. The virtual machine can load or store any of the control registers, enter or leave extended control mode, take interrupts, etc., without invoking the address translation feature.

If the virtual machine, already in extended control mode, turns on the translate bit in the EC mode PSW, then the routine DMKVATMD is called to examine the virtual control registers and build the required shadow tables. (Shadow tables are required since the real DAT hardware is capable of only a first-level storage mapping.) DMKVATMD examines virtual control registers 0 and 1 to determine if they contain valid information for use in constructing the shadow tables. Control register zero specifies the size of the page and

segment the virtual machine is using in the virtual page and segment tables. The shadow tables constructed by DMKVATMD are always in the same format as the virtual tables.

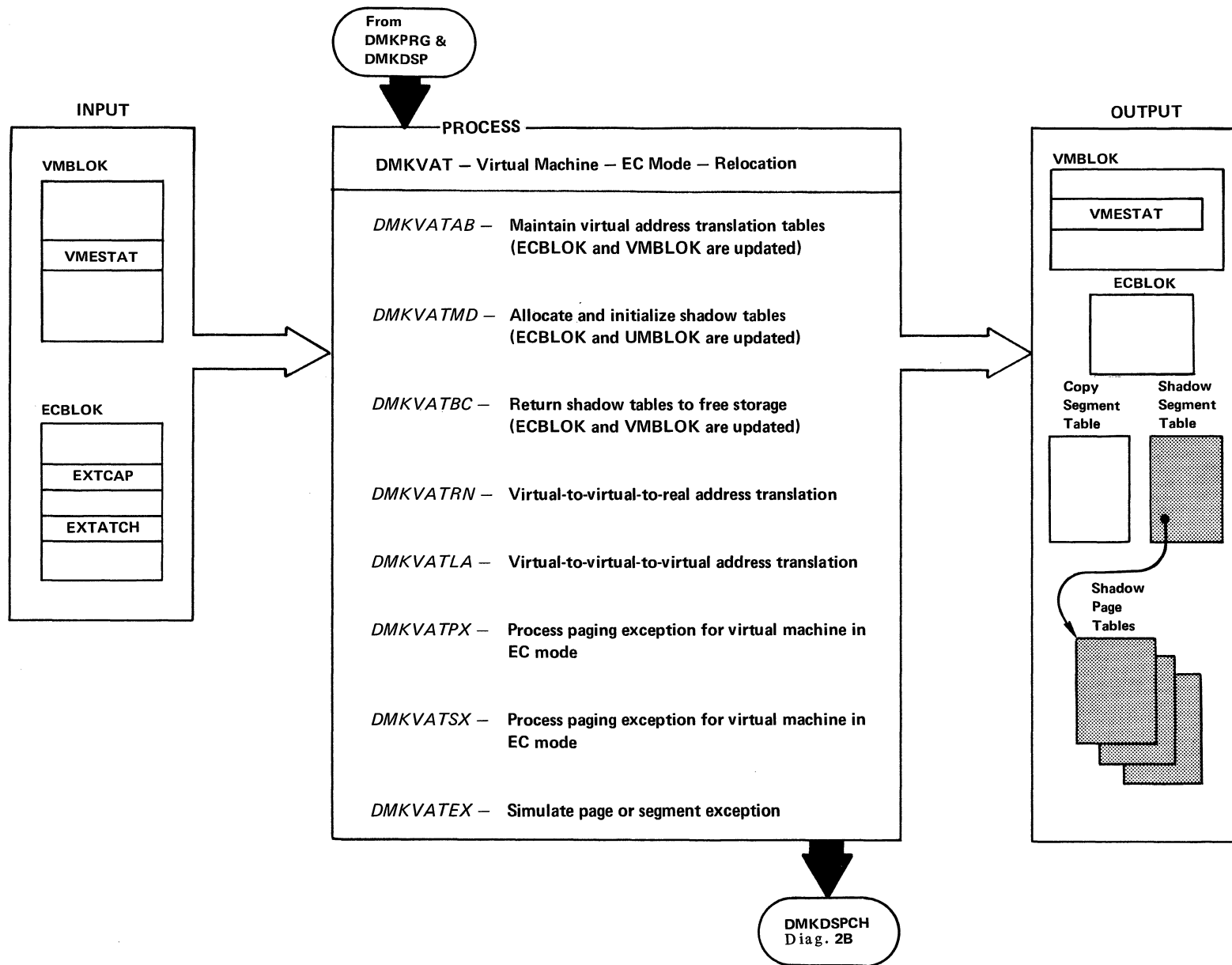
First, the virtual segment table is copied intact from second-level storage into first-level storage for speed of access when handling relocation interrupts. Another segment table of the same size, the shadow segment table, is constructed in first-level storage and initialized to indicate that all segments are unavailable. Flags are maintained in the VMBLOK to indicate that the shadow tables exist. DMKVATMD also constructs the shadow control registers 0 and 1. Shadow control register 0 contains the external interrupt mask bits used by CP, mixed with the hardware controls and enabling bits from virtual control register 0. Shadow control register 1 contains the segment table origin address of the shadow segment table.

When the virtual machine is operating in virtual translate mode, CP loads the shadow control registers into the real control registers and dispatches the user. The immediate result of attempting to execute an instruction is a segment exception, intercepted by DMKPRG and passed to DMKVATSX. DMKVATSX examines the copy, in first-level storage, of the virtual segment table in second-level storage. If the copy segment table indicates the segment is not available, the corresponding entry in the virtual segment table is examined and if necessary, the copy segment table is updated. If the virtual segment is not available, the segment exception interrupt is reflected to the virtual machine. If the virtual segment is marked available, then DMKVATSX:

- Allocates one full segment of shadow page table, in the format specified by virtual control register 0.
- Sets all of the page table entries to page not in storage.
- Marks the segment available in the shadow segment table.
- Redispatches the virtual machine via DMKDSP.

Once again, the immediate result is an interrupt, which this time is a paging exception and control is passed to DMKVATPX. DMKVATPX references the virtual page table

Diag. 1B3.6. Virtual Relocation



in second-level storage through the copy segment table to determine if the virtual page is available. If the virtual page is not available, the paging interrupt is reflected to the virtual machine. However, if the virtual page is marked in storage, the virtual page table entry is used to determine which page of second-level storage is being referenced by the third-level storage address provided. DMKVATPX next determines if that page of second-level storage is resident in first-level storage at that time. If so, the appropriate entry in the shadow page table is filled in and marked in storage. If not, the required page is brought into first level storage via DMKPTRAN and the shadow page table filled in as above.

As the virtual machine continues execution, more shadow tables are filled in or allocated as the third-level storage locations are referenced. Whenever a new segment is referenced, another segment of shadow page tables is allocated. Whenever a new page is referenced, the appropriate shadow pagetable entry is validated, etc. No changes are made in the shadow tables if the virtual machine leaves translate mode (usually via an interrupt), unless it also leaves extended control mode. Dropping out of EC mode is the signal for CP to release all of the shadow page and segment tables and the copy of the virtual segment table.

There are some situations that require invalidating all of the shadow tables constructed by CP or even releasing and reallocating them. Whenever DMKPTR swaps out a page that belongs to a virtual relocating machine, it sets a bit in the VMBLOK indicating that all of the shadow page tables must be invalidated. Invalidation of all of the tables is required since CP does not know which third-level-storage pages map into the second-level page which is being swapped out. The actual invalidation is handled by DMKVATAB, called from DMKDSP when the virtual machine is on the verge of being dispatched.

The other situations which cause shadow-table invalidation arise from the simulation of privileged instructions in DMKPRV. Flags are set in the VMBLOK whenever the virtual machine loads either control register 0 or 1, and DMKPRV calls DMKVATAB to perform whatever maintenance is required. When control register 1 is loaded by the virtual machine, DMKVATAB must recopy the virtual segment table into first-level

storage and invalidate the entire shadow segment table. When control register 0 is loaded, DMKVATAB examines the relocation-architecture control bits to determine if they have changed, (such that the format of the virtual page and segment tables no longer matches that of the shadow tables). If the format has not changed, the shadow tables are left intact; otherwise, all of the shadow tables and the copy segment table must be returned to free storage and another set, in the new format, must be allocated and initialized. The same actions can result from modifying the control registers via the CP console functions, in which case DMKVATAB is called from DMKCDB. The privileged operation, PTLB also causes the virtual segment tables to be recopied and all of the shadow page tables to be invalidated. since the shadow tables are the logical equivalent of the translation look-aside buffer.

DMKPRV provides virtual interrogation of the reference and change bits in the virtual storage keys, which involve the privileged instructions ISK, SSK, and RRB. The privileged instruction LRA is simulated via DMKVATLA, which searches the virtual page and segment tables to translate a third-level storage address to a second-level storage address, returning a condition-code indicator to DMKPRV, or forcing an interrupt if the tables are incorrectly formatted.

Most error situations that occur in the virtual machine are handled by means of the extended program interrupts associated with the real address translation hardware. Whenever a virtual relocating machine loads control registers 0 or 1 with an invalid value, DMKVAT releases all of the shadow tables and the copy segment table exactly as if the hardware controls had changed. The shadow control registers are set valid, with the shadow segment table re-allocated at a minimum size and all segments marked unavailable. Flag bits are set in the VMBLOK to indicate that the shadow tables are artificially valid, and DMKVATSX reflects a translation specification exception to the virtual machine as soon as it is dispatched. While it is possible for the virtual machine to enter an interrupt loop (if the new PSW is also a translate-mode PSW), the cited process prevents the occurrence of a disabled-loop within CP, which would result if the virtual machine is never dispatched.

PRIVILEGED INSTRUCTIONS

If the program interruption is caused by the virtual machine issuing a privileged instruction, DMKPRVLG obtains the address of the privileged instruction and determines the type of operation requested.

I/O Privileged Instructions

DMKPRVLG transfers control to the virtual I/O executive program (DMKVIOEX).

Non-I/O Privileged Instructions

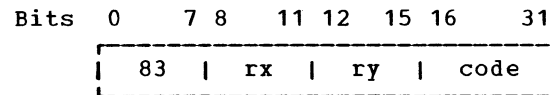
DMKPRVLG simulates valid non-I/O privileged instructions and returns control to DMKDSPCH. For invalid privileged instructions, the routine sets an invalid interruption code and reflects the interruption to the virtual machine. For the privileged instructions SCK, SCKC, STCKC, SPT, and STPT that affect the TOD clock, CPU timer, and TOD clock comparator, control is transferred to DMKTMR by DMKPRVLG. Others that are simulated are LPSW, SSM, SSK, ISK, and diagnose.

System/370 EC mode privileged simulation includes the following:

<u>Code</u>	<u>Definition</u>
SCK	Set clock
SCKC	Set clock comparator
STCKC	Store clock comparator
SPT	Set CPU timer
STPT	Store CPU timer
STNSM	Store and AND system mask
STOSM	Store and OR system mask
STIDP	Store CPU identification
STIDC	Store channel identification
LCTL	Load control
STCTL	Store control
LRA	Load real address
RRB	Reset reference bit
PTLB	Purge table look-aside buffer

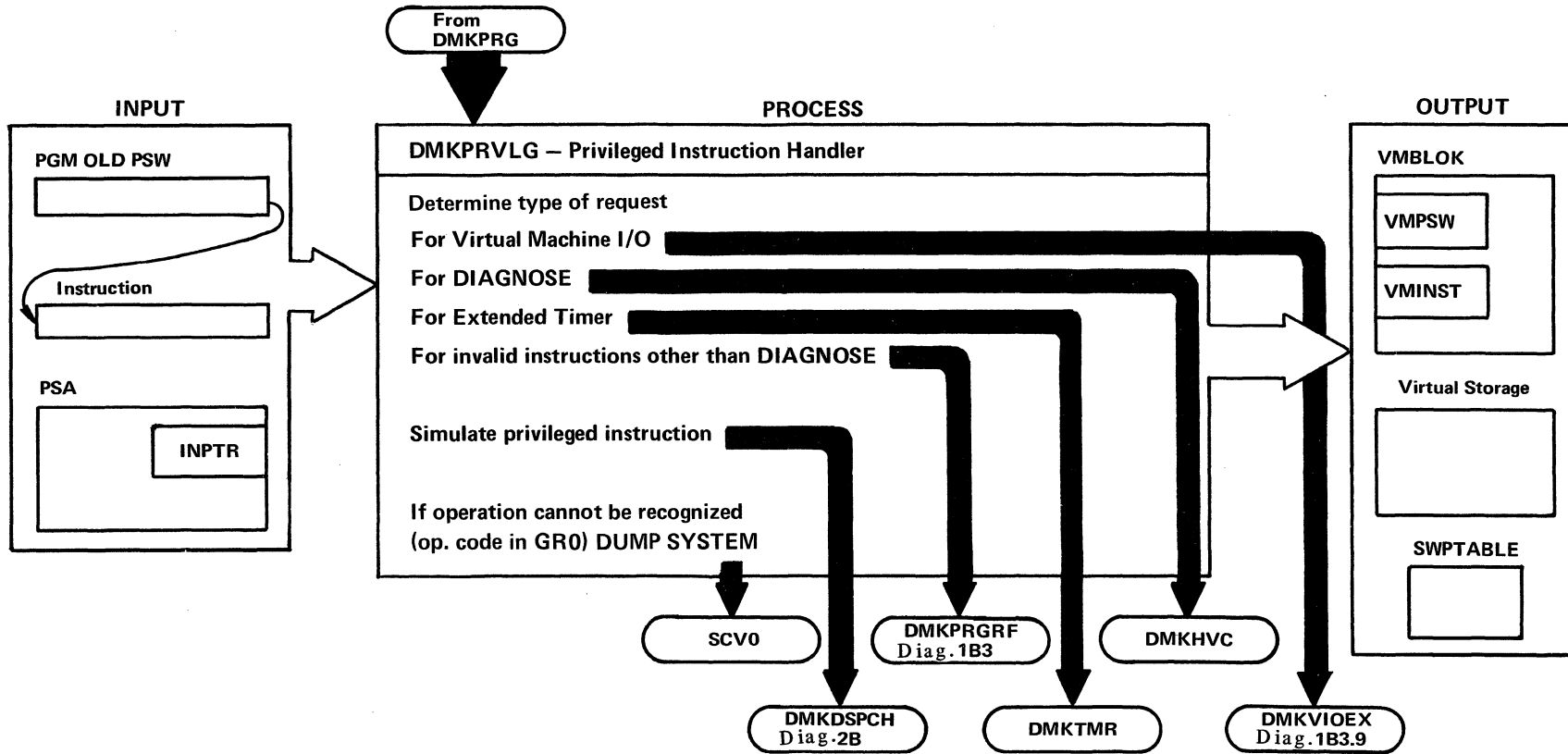
DIAGNOSE Interface (DMKHVC)

The diagnose command is used for communication between a virtual machine and the VM/370 control program. In VM/370, the machine-coded format for the diagnose command is:



- 83 is the Diagnose operation code.
- rx is a user specified register number.
- ry is a user specified register number.

Diag. 1B3.7. Privileged Instruction Simulation



Code is a hexadecimal value that is used to select a particular VM/370 control program function. The codes and their associated functions are:

Code	Class	Function
0004	C,E	Examine data from real storage
0008	G	Execute VM/370 control program console function
000C	G	Pseudo-timer facility
0010	G	Release virtual storage pages
0014	G	Manipulate input spool files
0018	G	Standard DASD I/O
001C	F	Clear I/O and machine check recording
0020	G	General virtual I/O without interrupts
0024	G	Virtual device type information
002C	C,E,F	Return DASD start of LOGREC area
0030	C,E,F	Read one page of LOGREC data
0034	C,F	Read system-dump spool file
0038	C,E	Read system symbol table
003C	A,B,C	Dynamically update system user directory
004C	Any	Generate accounting cards for virtual user

Notes: Rules for diagnose codes:

X'00' through X'FC' Reserved for IBM use.  
 X'100' through X'1FC' Reserved for users

The diagnose code must always be a multiple of 4.

DIAGNOSE CODE 4: Examine real storage, can only be issued by users with privilege class C or E.

rx contains the virtual address of a list of CP (real) addresses.

ry (cannot be register 15) contains a count of entries in the list.

ry+1 contains the virtual address of the result field that holds the values retrieved from the VM/370 control program locations.

DIAGNOSE CODE 8: Virtual console function, allows a virtual machine to perform the VM/370 control program console functions.

rx contains the address (virtual) of the control program console function command and parameters.

ry contains the length of the associated console function input, up to 132 characters.

The following illustrates the virtual console function:

```

LA R6,CPFUNC
LA R10,CPFUNCL
DC X'83',X'6A',XL2'0008'
.
.
CPFUNC DC C'QUERY FILES'
CPFUNCL EQU *-CPFUNC
  
```

The output of the console function is to the user's terminal, and then execution continues. Any valid and authorized console function can be executed in this manner.

A completion code is returned to the user as a value in the register specified in ry. The error code = the message number of the error message issued.

DIAGNOSE CODE C: Pseudo timer.

ix contains the virtual address of a 32-byte data area that does not cross a page boundary, into which the following data is stored:

Bytes	0	7	8	15	16	23	24	31
	MM/DD/YY		HH:MM:SS		Virt CPU		Total CPU	

Virtual and Total CPU time used is returned as a doubleword logical value in microseconds.

DIAGNOSE CODE 10: Release pages.

ix contains the virtual address of the first page to be released.

ry contains the virtual address of the last page to be released.

Any of the virtual pages in real or auxiliary storage are released.

DIAGNOSE CODE 14: Input spool file manipulation.

rx contains either a buffer address, a copy count, or a spool-file identifier, dependent on the value of the function subcode in ry+1.

ry (cannot be register 15) contains the virtual address of a spool-input card reader.

ry+1 contains a function hexadecimal code interpreted by DMKDRDER, as follows:

<u>Code</u>	<u>Function</u>
0000	Read next spool buffer (data record)
0004	Read next print SFBLOK
0008	Read next punch SFBLOK
000C	Select a file for processing
0010	Repeat active file <u>nn</u> times
0014	Restart active file at beginning
0018	Backspace one record

ry+1 on return, may contain error codes which further define a returned condition code of 3. See Figure 1 for Condition Code analysis.

The file manipulation is performed by DMKDRDER.

DIAGNOSE CODE 18: Disk I/O.

rx contains the device address of the disk.

ry points to a CCW chain to read or write a limited number of disk records.

Each read or write must specify no more than 2048 bytes (usually 800 is used), and the CCW chain is of a standard form, as shown below. For a 3330, a SET SECTOR command would precede each SRCH command.

Register 15 contains the number of reads or writes in the CCW chain (the number is two in the following example for a typical CCW string (to read or write two 800-byte records):

```

SEEK,A,CC,6
SRCH,A+2,CC,5
TIC,*-8,0,0
RD or WRT,DATA,CC+SILI,800
SEEK HEAD,B,CC,6 (Omitted if HEAD No. unchanged)
SRCH,B+2,CC,5
TIC,*-8,0,0
RD or WRT,DATA+800,SILI,800
A  SEEK and SRCH arguments for first RD/WRT
B  SEEK and SRCH arguments for second RD/WRT

```

DIAGNOSE CODE 1C: Clear I/O recording, can only be issued by a privilege class F user. This code calls the DMKIOEFM routine to clear the I/O error recording data on disk.

rx contains the code value 1, 2, or 3 to clear and reformat the I/O error recording, M/C recording, or both I/O and M/C recording, respectively.

ry is ignored.

DIAGNOSE CODE 20: General I/O without interrupts.

rx contains a virtual device address.

ry contains the address of the string of CCWs to be executed.

The CCW string is processed via DMKCCWTR through DMKGENIO, providing full virtual I/O in a synchronous fashion (self-modifying CCW strings are not permitted, however) to any virtual device specified. Control returns to the virtual machine only after completion of the operation or detection of a fatal error condition. Condition codes and error codes in ry are returned to the virtual system.

DIAGNOSE CODE 24: Virtual device type information.

rx contains a virtual device address.

ry, which cannot be register 15, and ry+1 contain the following upon return:



Code	ry+1	Meaning
0		Data Transfer successful
1		End of file
2		File not found
3	4	Device address invalid
3	8	Device type invalid
3	12	Device busy
3	16	Fatal Paging I/O Error

Figure 1. Condition Code Analysis for Diagnose Codes 14 and 34

Bits	0	7 8	15 16	23 24	31
ry	VDEVTYPEPC VDEVTYPE VDEVSTAT VDEVFLAG				
ry+1	RDEVTYPEPC RDEVTYPE RDEVMDL  RDEVFIR				

A condition code of one or three indicates that the virtual device address specified is either invalid (that is, too large), or the device does not exist. Condition code 2 indicates the real device does not exist.

**DIAGNOSE CODE 2C:** Return DASD start of LOGREC area (Privilege class C, E, or F only).

rx on return contains the DASD location, in VM/370 control program internal format, of the first record of the system I/O and machine check error recording area.

ry is ignored.

**DIAGNOSE CODE 30:** Read one page of LOGREC data (Privilege class C, E, or F only).

rx contains the DASD location, in the VM/370 control program internal format, of the desired record.

ry contains the virtual address of a page-size buffer to receive the data.

The page of data is provided to the virtual machine via DMKRPAGT.

cc = 0 Successful read, data available.  
 1 End of cylinder, no data.  
 2 Invalid cylinder, outside recording area

**DIAGNOSE CODE 34:** Read system dump spool file (privilege class C or E only).

rx contains the virtual address of a page-size buffer to accept the requested data.

ry (cannot be register 15) contains the virtual device address of a spool-input card reader.

ry+1 on return, may contain error codes which further define a returned condition code of 3. See Figure 1 for Condition Code analysis.

The system chain of spool input files is searched for a dump file belonging to the user issuing the diagnose command by DMKDRDMP. The first (or next) record from the dump file is provided to the virtual machine via DMKRPAGT and the condition code is set to zero. The dump file is closed via VM/370 console function CLOSE.

**DIAGNOSE CODE 38:** Read system symbol table.

rx contains the start address of the page buffer that is to contain the symbol table.

ry is ignored.

The system symbol table (DMKSYM) is read into storage at the location specified by rx by DMKDRDSY.

**DIAGNOSE CODE 3C:** Dynamically update the system user directory.

rx contains the first 4 bytes of the volume serial label.

ry, the first 2 bytes of the register specified (ry) contain the last 2 bytes of the volume serial label.

The directory is dynamically updated by DMKUDRDS.

DIAGNOSE CODE 4C: Generate accounting cards for the virtual user. This code can be issued only by a user with the account option (ACCT) in his directory.

rx contains the virtual address of a 24-byte parameter list identifying the "charge to" user; the address must be aligned on a doubleword boundary. If rx contains zeros, the accounting card will be punched with the identification of the user issuing the diagnose instruction.

ry contains a function hexadecimal code interpreted by DMKHVC as follows:

Code Function

0000 The parameter list contains only a userid.  
 0004 The parameter list contains a userid and account number.  
 0008 The parameter list contains a userid and distribution number.  
 000C The parameter list contains a userid, account number, and distribution number.

The following condition codes are returned to the user by DMKHVC:

cc=0 Successful operation  
 1 User does not have account option privileges  
 2 Invalid userid in the parameter list  
 3 Invalid function hexadecimal code in ry or an error occurred in trying to read in the User Machine Block (UMACBLOK)

DMKHVC checks that the user has the account option and if not, returns a condition code of 1. If the user has the options, control is passed to DMKCPV to generate the card. DMKCPV passes control to DMKACO to complete the "charge to" information; either from the User Accounting Block (ACCTBLOK), if a pointer to it exists, or from the user's VMBLOK. DMKCPV then punches the card and passes control back to DMKHVC to release the storage for the ACCTBLOK, if one exists. DMKHVC then checks the parameter list address for the following conditions:

- If zero, control is returned to the user with a condition code of zero.

- If invalid, an addressing exception is generated.
- If not aligned on a doubleword boundary, a specification exception is generated.

For a parameter list address that is non-zero and valid, the userid in the parameter list is checked against the directory list and if not found, control is returned to the user with a condition code of two. If the function hexadecimal code is invalid, control is returned to the user with a condition code of three. If both userid and function hexadecimal code are valid, the User Accounting Block (ACCTBLOK) is built and the userid, account number, and distribution number are moved to the block from the parameter list or the User Machine Block belonging to the userid in the parameter list. Control is then passed to the user with a condition code of zero.

Virtual Timer Maintenance

The System/370 with EC mode provides the system user (both real and virtual) with four timing facilities. They are:

1. The interval timer at main storage location X'50'.
2. The time-of-day clock.
3. The time-of-day clock comparator.
4. The CPU timer.

REAL TIMING FACILITIES: Before describing how CP maintains these timers for virtual machines, it is necessary to review how VM/370 uses the timing facilities of the real machine.

1. The location X'50' interval timer is used only for time-slicing. The value placed in the timer is the maximum length of time that the dispatched user is allowed to execute.
2. The time-of-day clock is used as a time stamp for messages and enables the scheduler to compute

elapsed in-queue time for the dispatching priority calculation.

3. The time-of-day clock comparator facility is used by CP to schedule timer driven events for both control program functions and for virtual machines. A stack of comparator requests is maintained and as clock comparator interrupts occur, the timer request blocks are stacked for the dispatcher via calls to DMKSTKIO.
4. The CPU timer facility performs three functions:
  - Accumulation of CP overhead
  - Detection of in-queue time slice end
  - Virtual CPU timer simulation

The accumulation of CP overhead is accomplished as follows. The VMTIME field in the VMBLOK contains the total CP overhead incurred by the virtual machine; it is initialized to the maximum sized doubleword integer, X'7FFFFFFF FFFFFFFF'. Whenever CP is to perform a service for a virtual machine, GPR 11 is loaded with the address of the VMBLOK and the current value in VMTIME is placed in the CPU timer. When CP is finished with the service for that virtual machine the CPU timer, which has been decremented by the amount of CPU time used, is stored back into VMTIME. GPR11 is then loaded with a new VMBLOK pointer and the CPU timer is set from the new VMTIME field. The amount of CP overhead for a given virtual machine at any point in time is the difference between the maximum integer and the current value in the VMTIME field.

Since VMTIME only accounts for supervisor state overhead, detection of in-queue time slice end is performed by the CPU timer when the virtual machine is dispatched in the problem state. The VMTMOUTQ field in the VMBLOK is initialized to the amount of problem state time that the virtual machine will be allowed to accumulate before being dropped from a queue. This initial value is set by the scheduler (DMKSCH) when the virtual machine is added to a queue and its value depends on the queue entered (interactive or non-interactive) and on the CPU model. For example, the initial value of VMTMOUTQ for a user entering Q1 (interactive)

on a model 145 is 300 milliseconds, while for the same user entering Q2 (non-interactive) it is 2 seconds. Each time the user is dispatched, the value in VMTMOUTQ is entered into the CPU timer; whenever the user is interrupted, the decremented CPU timer is stored into VMTMOUTQ prior to being set from the new VMTIME. When the problem state time slice has been exhausted; a CPU timer interrupt occurs, the VMQSEND flag bit is set in the VMBLOK, and the scheduler drops the user from the queue. At each queue drop, the problem time used in-queue (the difference between VMTMOUTQ and the initial value) is added to the total problem time field (VMVTIME) in the VMBLOK.

Virtual CPU timer simulation is handled for EC mode virtual machines if the value in their virtual CPU timer is less than that in VMTMOUTQ. In this case, the VMBLOK is flagged as "tracking CPU timer" and a CPU timer interrupt is interpreted as a virtual timer interrupt rather than as an in-queue time slice end.

VIRTUAL TIMING FACILITIES: Virtual location X'50' timers are updated by the elapsed CPU time each time the dispatcher has been entered after a running user has been interrupted. The size of the update is the difference between the value of the timer at dispatch (saved in QUANTUM at location X'54') and the value of the timer at the time of the interrupt (saved in QUANTUMR at location X'4C').

Virtual clock comparator requests are handled by the virtual timer maintenance routine DMKTMR. They are inserted into the general comparator request stack and the virtual machine is posted when the interrupt goes off.

Requests to set the virtual CPU timer place the new value into the ECBLOK. Requests to store it update the ECBLOK field by the virtual CPU time used since the last entry to dispatch and pass the value to the user. Requests to set the time of day clock are ignored.

A real interval or CPU timer is one which runs when the user is executing or is in a self-imposed wait state (that is, the wait bit is on in his virtual PSW). A real timer does not run if the user is in a CP pseudo-wait (for example, page wait or I/O wait) or if

he can be run but is not being dispatched due to other user interaction. Real timers provide accurate interrupts to programs that depend on measurement of elapsed CPU and/or wait time. They do not accurately measure wall time -- the TOD clock must be used for this function.

An EC mode virtual machine with the Real Timer option has both a real interval timer and a real CPU timer. Real timer requests for waiting machines are maintained in the clock comparator stack. CPU timer requests are added to TOD clock value at the time that they are issued. Interval timer requests must have their units converted. In addition, if the virtual CPU timer contains a large negative value, then a real timer request is scheduled to occur when the virtual time turns positive, so that the pending timer interrupt can be unflagged. Comparator requests for real timer interrupts are inserted into the stack whenever a user enters a self-imposed wait. They are removed either when the user resumes execution or when he is forced (or places himself) into a pseudo wait.

### Virtual I/O Requests

The function of the virtual I/O interface maintained by the control program is to provide to the software operating in the user's virtual machine the condition codes, CSW status information, and interrupts necessary to make it appear to the user software that it is in fact running on a real System/370. The virtual I/O interface consists of:

- A virtual I/O configuration of each active user represented by a set of I/O control blocks that are maintained in the Control Program's free storage. This configuration is built at LOGON time from information contained in the user's directory file, and can be changed by the user or the system operator.
- A set of routines that maintain in these blocks, the status of the virtual I/O configuration.
- Other system components to simulate/translate the channel programs provided by the user to initiate I/O on units in the real system's configuration.

VIRTUAL I/O CONTROL BLOCKS: The base for locating the I/O block structure is the user's Virtual Machine Block (VMBLOCK). The VMBLOCK contains a pointer to the start of three control block tables, and a table of 16 channel indexes. The control block tables contain one block for each of the virtual channels, control units, and devices that are defined for the user's virtual machine. The entries in the channel index table (VMCHTBL) contain the pointers to each channel defined for the user in the table of Virtual Channel Blocks (VCHBLOCKs). Each VCHBLOCK contains a table of pointers that point to the Virtual Control Unit Blocks (VCUBLOCKs) for the control units attached to that virtual channel. Each VCUBLOCK contains pointers to the Virtual Device Blocks (VDEVBLOCKs) attached to the control unit. See Diag. 1B3.8 for an overview of the virtual I/O control blocks.

Thus, if given the unit address of any component in the form ccu, the appropriate control blocks representing each component in the subchannel path to the given unit is located via the indexing scheme.

VCHBLOCK: There is one VCHBLOCK for each virtual channel connected to the user's virtual CPU. Each VCHBLOCK contains the channel address and flag indicating the channel type (selector, byte multiplexer or block multiplexer). The status of the channel and its attached units are represented by several status and mask bytes, as follows:

1. A status byte (VCHSTAT) indicates whether the channel is busy or has a channel class interrupt pending.
2. A halfword unit address identifies the unit causing the channel-class interrupt (if it is present).
3. A halfword mask (VCHCUINT) contains a bit map of the attached control units that have interrupt status pending. Following these status flags and masks is the table of indexes pointing to the attached VCUBLOCKs; index entries representing addresses at which no control unit is attached have a value of -1.

VCUBLOCK: There is one VCUBLOCK for each control unit in the virtual configuration. These blocks are arranged

in a table, and each contains, in addition to its base address, status flags similar to those in the VCHBLOK and a table of indexes to attached VDEVBLOKS. The status flags defined for the VCUBLOK differ from those for the VCHBLOK in that they can contain status for the control unit and also for a subchannel.

For example, if the VCUBLOK representing a 2803 Tape Control Unit is attached to a virtual selector channel, both the VCHBLOK and the VCUBLOK are marked busy. However, if the VCUBLOK is attached to a virtual byte multiplexer channel and is for a control unit on a selector subchannel of the multiplexer, the busy status of the channel is reflected in the VCUBLOK only. Thus the virtual multiplexer appears nonbusy to operations on other, nonshared subchannels.

VDEVBLOK: There is one VDEVBLOK in the configuration for each virtual device defined by the user. Each

VDEVBLOK contains the device portion of the unit address, device status, and the virtual CSW for the last interrupt taken by the device. In addition, the VDEVBLOK contains device type specific information that allows the I/O translation and simulation routines to interpret the channel programs presented by the user. This information is not used by the I/O interface.

Since all virtual machines are run in the problem state, any attempt to issue a SIO instruction results in a program interrupt that indicates a privileged operation exception. This interrupt is handled by CP's first level program interrupt handler, DMKPRGIN. It determines if the virtual machine was in virtual supervisor state (problem state bit in the VIRTUAL PSW is zero). If so, the instruction causing the interrupt is saved in the VMBLOK for the virtual machine and control is transferred to the privileged instruction simulator, DMKPRVLG, via a GOTO.

**Diag. 1B3.8 Virtual I/O Control Blocks**

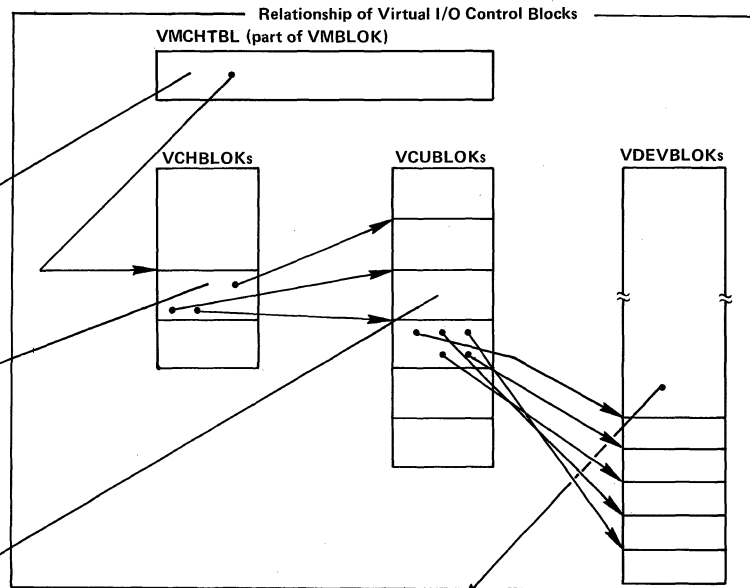
The virtual machine configuration is represented by a set of related control blocks. These blocks are:

- built by VM/370 at LOGIN from data in directory
- modified by user commands (e.g. DETACH, LINK, DEFINE)

There is one control block per channel, per control unit, and per device.

The characteristics of VM/370 virtual I/O control are:

- RSP (Rotational Position Sensing) cannot be used on BMPX (Block Multiplexer Channel)
- No multi-path configurations
- The virtual machine operating system performs scheduling
- VM/370 uses virtual I/O control blocks to simulate real hardware interface.
- Virtual unit record devices use VM/370 spooling
- Virtual console is simulated on terminal
- Mini-disks simulate DASD
- Dedicated devices are supported.



VMCHTBL – virtual channel index table

VCHBLOK – virtual channel block

Channel Identification Status			
XXXX	XXXX	XXXX	XXXX
XXXX	XXXX	XXXX	XXXX
XXXX			

XXXX

If negative (FFFF), no control unit exists  
 If negative (8XXX) the control unit exists but the VCUBLOK cannot be addressed by the virtual machine because the control unit is detached.  
 If positive, the value is on index to the VCUBLOK.

VCUBLOK – virtual control unit block

Control unit identification status			
XXXX	XXXX	XXXX	XXXX
XXXX			
XXXX			

Device Index Table

XXXX

If negative (FFFF), no device exists  
 If negative (8XXX) the device exists but the VDEVBLOK cannot be addressed by the virtual machine because the device is detached.  
 If positive, the value is an index to the VDEVBLOK.

VDEVBLOK – virtual device block

Device identification
Status pending
Positioning
Terminal control
Spooling control
...
RDEVBLOK pointer

Part of the VDEVBLOK contains device independent information and is used identically in all VDEVBLOKs. However, some fields of the VDEVBLOKs have multiple uses, depending on the device type.

DMKPRVLG determines if the privileged operation affects the virtual I/O configuration. DMKPRVLG simulates non-I/O privileged instructions (such as LPSW) itself. If the instruction's operation code is from X'9C to X'9F', control is transferred to DMKVIOEX.

After clearing the condition code in the user's VMBLOK, DMKSCNVU is then called to locate the Virtual I/O blocks representing the components (channel, control unit and device) addressed by the instruction. DMKVIOEX then branches to handle the request based on the operation requested.

VIRTUAL SIO (See Figure 2): With a SIO, the condition code returned from DMKSCNVU is tested to verify that all addressed components were located. If they were not, then a condition code of 3 (unit not available) is reflected in the VPSW and control is returned to the dispatcher. Otherwise, the addresses of the appropriate virtual I/O control blocks are saved, and DMKVIOEX tests the status of the addressed I/O units by scanning the VCHBLOKS, VCUBLOKS, and VDEVBLOKS to locate the block that contains the status of the addressed subchannel. The subchannel status is indicated in:

- The VCHBLOK for a selector or block multiplexer channel
- The VCUBLOK for a shared selector subchannel on a byte multiplexer.
- The VDEVBLOK for a nonshared subchannel on a byte multiplexer.

When the block containing the status is found, the status is tested. If the subchannel is busy or has an interrupt pending, condition code 2 is reflected. Otherwise, the subchannel is available and the device and the control unit are tested for interrupt pending or busy. If either is found, condition code 1 is reflected and the proper CSW status is stored in the user's virtual page zero. If all components in the subchannel path are free, DMKVIOEX proceeds to simulate the SIO by locating and loading the contents of the user's CAW from his virtual location X'48' and testing the device type of the unit addressed.

The device type is determined by referencing the VDEVBLCK. If the device class code indicates a terminal or console, control is passed to the virtual

console executive DMKVCNEX via a GOTO. DMKVCNEX interprets and simulates the entire channel program, moving the necessary data to or from the user's virtual storage and reflecting the proper interrupts and status bytes. When DMKVCNEX has finished, it passes control directly to the dispatcher DMKDSPCH.

If DMKVIOEX determines that the referenced device is a spooled unit-record device, it passes control to DMKVSPEX for additional processing and upon return it passes control to DMKDSPCH.

If the referenced device is not a terminal nor a spooling device, the SIO is translated and executed directly on the real system's I/O device. DMKVIOEX calls DMKFREE to obtain free storage and then it constructs an IOBLOK in the storage obtained. The IOBLOK serves as an identifier of the I/O task to be performed. It contains a pointer to the channel program to be executed and the address of the routine that is to handle any interrupts associated with the operation.

DMKVIOEX stores the contents of the user's CAW in IOBCAW and sets the interrupt return address (IOBIRA) to the virtual interrupt return address (DMKVIOIN) in DMKVIO. The CCW translation routine (DMKCCWTR) is then called to locate and bring into real main storage all user pages associated with the channel program, including those containing data and CCWs.

1. The CCWs are translated.
2. A corresponding real channel program is constructed.
3. The data pages are locked into real storage.
4. DMKCCWTR returns control to DMKVIOEX. DMKVIOEX places the user in a pseudo-wait state, IOWAIT, and calls the real I/O scheduler DMKIOSQV to schedule the I/O on the real configuration.

DMKIOSQV queues the request for operation on the real channel, control unit, and device corresponding to the one addressed by the user. When the real SIO is issued, DMKIOS takes the user out of IOWAIT and reflects the condition code for the SIO if it is zero. If it is not zero, the operation is further analyzed by DMKVIOIN. In any case, DMKIOSQV returns control to DMKVIOEX, which passes control to DMKDSPCH.

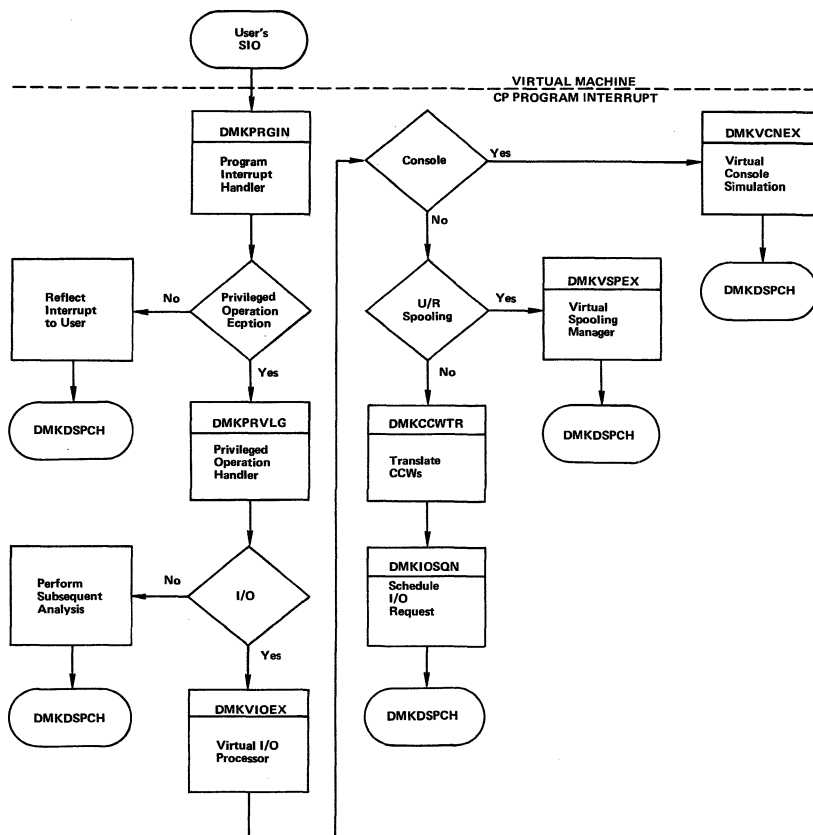


Figure 2. Overview of a Virtual SIO

OTHER VIRTUAL I/O INSTRUCTIONS: Other privileged I/O instructions are handled directly by DMKVIOEX. The general method used is to scan the virtual channel, control unit, and device blocks in the same manner as for the SIO and to reflect the proper status and condition to the user. In some cases (TIO), the status of the addressed components are altered after the status is presented.

If the operation active on the virtual device is actually in progress in the real equipment, the simulation of a HIO or HDV is somewhat more involved, since it requires the actual execution of the instruction. In this case, the active operation is halted and the resultant condition code/status is returned to the user.

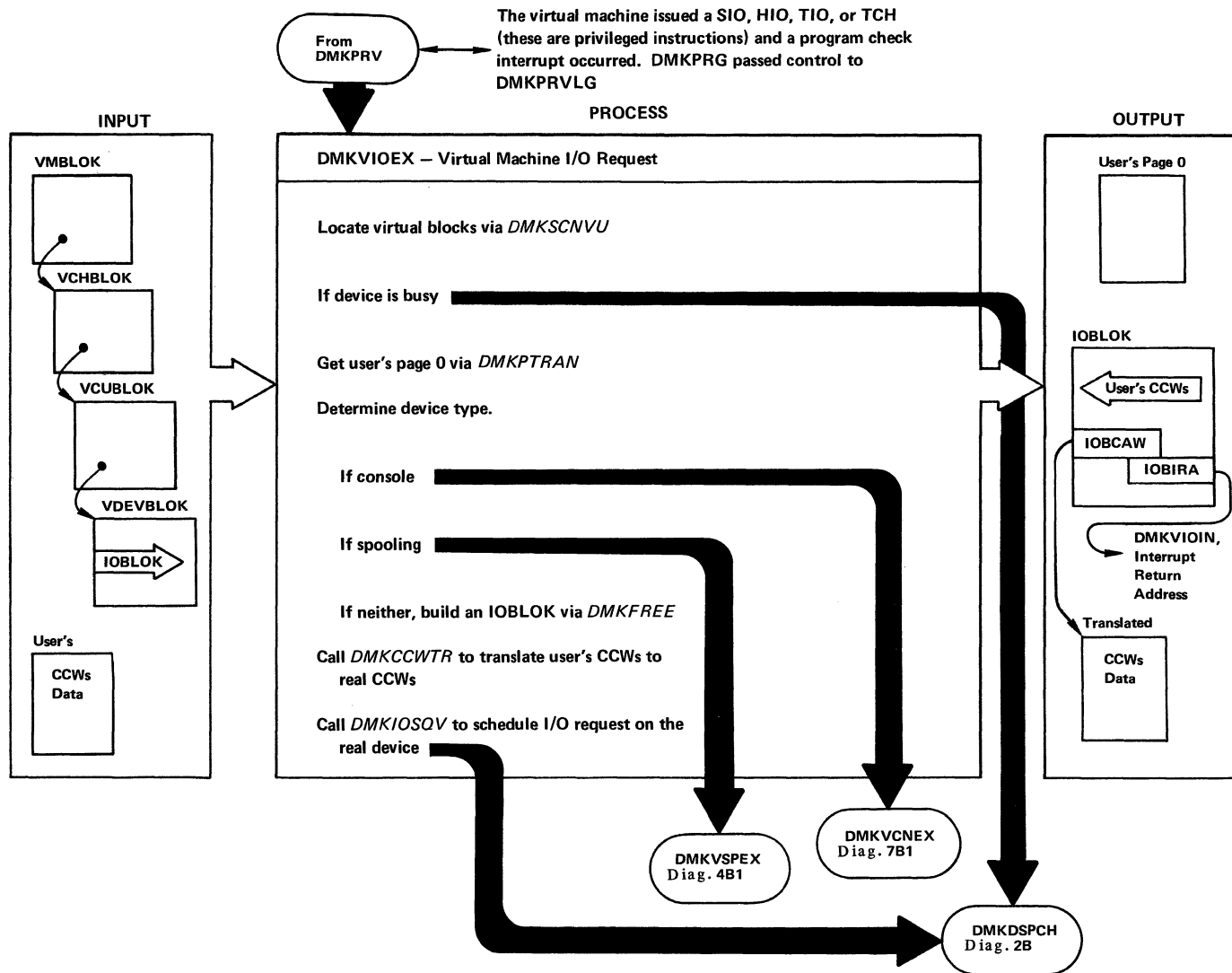
VIRTUAL CHANNEL-TO-CHANNEL ADAPTER: The virtual channel-to-channel adapter (CTCA) is simulation that permits data transfer and control communication between two selector channels, either on two distinct processors or two channels on a single processor. Data transfer is accomplished via synchronized complementary I/O commands (for example, read/write, write/read) issued to both parts of the CTCA. Each part of the CTCA is identical and the operation of the unit is completely symmetrical. The CTCA occupies an entire control-unit slot on each of the two channels attached. The low-order four bits of the unit address (device address) are ignored completely and are not available for use.

The VM/370 control program support for virtual CTCA includes all status, sense data, and interrupt presentation logic necessary to simulate the operation of the real CTCA. Data transfer, command byte exchange, sense data, and status data presentation for the virtual CTCA is accomplished via storage-to-storage operations (MVCL, etc.). No real I/O operations (excluding paging I/O) nor I/O interrupts are involved. Unit errors or control errors cannot occur.

VIRTUAL SELECTOR CHANNEL I/O REQUESTS: The CCW translator, DMKCCWTR, is called by the virtual machine I/O executive program (DMKVIOEX) when an I/O task block has been created and a list of virtual CCWs associated with a user's SIO request must be translated into real CCWs.



Diag. 1B3.9. Virtual I/O Request



When the I/O operation from a self-modifying channel program is completed, DMKUNTIS is called by DMKIOS. When retranslation of OS ISAM CCWs is required, the self-modifying channel program checking portion of DMKCCWTR calls DMKISMTR.

DMKCCWTR operates in two phases:

- A scan and a translate phase.
- A TIC-scan phase, if the ISAM option was chosen.

A self-modifying channel program checking function is also included.

The scan and translate phase analyzes the virtual CCW list. Some channel commands require additional doublewords for control information (for example, seek addresses). Additional control words are also allotted (in pairs) if the data area specified by a virtual CCW crosses 4096-byte page boundaries, or if the virtual CCW includes an IDA (Indirect Data Address) flag.

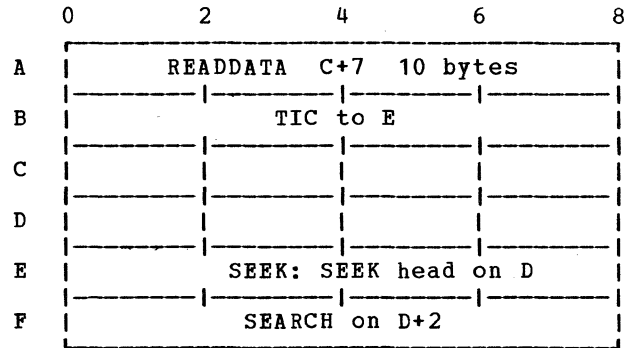
Space is obtained from DMKFREE for the real CCW list, and the translation phase then translates the virtual CCW list into a real CCW list. TIC commands that cannot be immediately translated are flagged for later processing by the TIC-scan phase. A read or write command that specifies data crossing 4096-byte boundaries is revised to include an IDA flag that points to an Indirect Data Address List (IDAL) and a pair of words for each 4096-byte page, in which each word handles a data-transfer of 2048 bytes (or less). The real CCW is flagged as having a CP-generated IDA. DMKPTRAN is called (via the TRANS macro) to lock each 4096-byte page.

If the real CCW string does not fit in the allocated free storage block, a new block is obtained. The old block is transferred and adjusted before being released. The translation continues with the new block. The process is repeated as needed to contain the real CCW string.

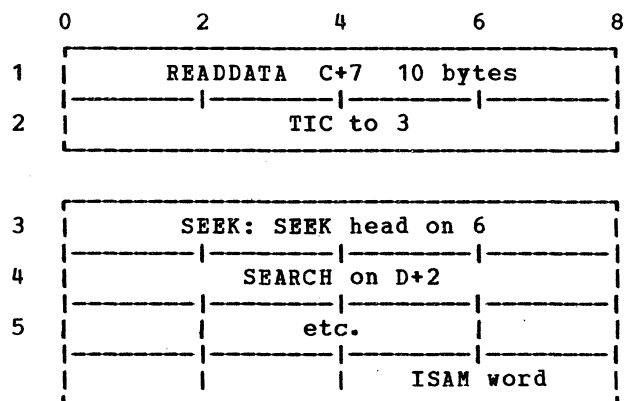
Virtual CCWs having and IDA flag set are converted to use translated addresses for each IDAW (Indirect Data Address Word) in the virtual IDAL. DMKPTRAN is called for each IDAW. The CCW flagged as having a user (but not CP) generated IDA.

The TIC-scan phase scans the real CCW list for flagged (untranslated) TIC commands and creates a new virtual CCW list for the untranslated commands. Scan-translate phase processing is then repeated. When all virtual CCWs are translated, the virtual CAW in the IOBLOK task block is replaced by the real CAW (that is, a pointer to the real CCW list created by DMKCCWTR), and DMKCCWTR returns control to DMKVIOEX. The user protection key is preserved.

OS ISAM Handling by DMKISMTR: Because many of the OS PCP, MFT, and MVT ISAM channel programs are self-modifying, special handling is required by the VM/370 control program to allow virtual machines to use this access method. The particular CCWs that require special handling have the following general format:



The CCW at A reads 10 bytes of data, the last byte of which forms the command code of the CCW at E. In addition, the data read in forms the seek and search arguments for the CCWs at E and F. After the CCW string is translated by the VM/370 control program it usually is in the following format:



In order to accomplish an efficient and non-timing dependent translated operation for OS ISAM, the virtual CCW string is modified in the following manner.

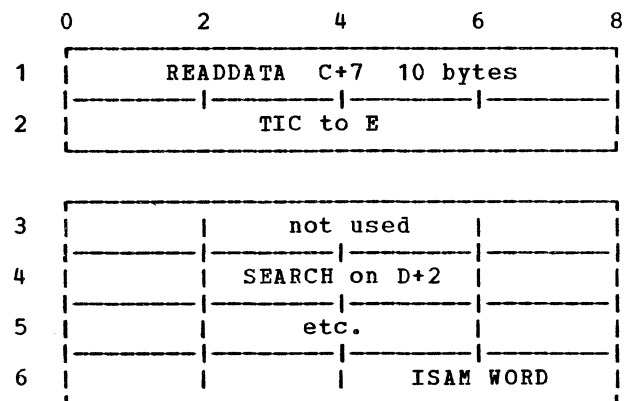
DMKISMTR is called by DMKCCWTR if, during normal translation, a CCW of the type at 1 is encountered. The scan program locates the TIC at 2 by searching the translated CCW strings. The TIC at 2 locates the seek at 3.

The virtual address of the virtual seek CCW at E is located from the RCWTASK header. Three doublewords of free storage are obtained and the address of the block is saved in the ISAM control word at 5. The three doublewords are used to save the following information from the translated CCW strings and from the users virtual storage.

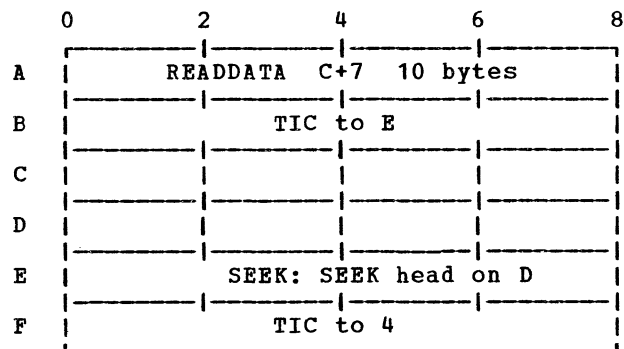
Before	After
address of TIC at 2	first word of TIC at 2
address of SEEK at E	first word of SEEK at E
first word of CCW at F	address of CCW at 4

The TIC at 2 is altered to TIC to the virtual CCW at E. The CCW address field at E is translated to reference D. The four bytes at F are modified to a TIC to the

CCWs starting at 4. The completed CCW string has the following format:



#### TRANSLATED CCWs



This interrupt return address in the IOBLOK is set to DMKUNTIS. DMKUNTIS restores the data to its original format from the three doubleword extension and releases the block. Normal I/O handling is resumed by DMKVIO and DMKUNT.

#### I/O SUPERVISOR

The module DMKIOS handles the I/O requirements of all system devices except for the low-speed lines that

serve as user logon consoles. Scheduling and interrupt supervision for these devices is essentially a synchronous process and does not require the queuing and restart services of DMKIOS; it is therefore handled by the module DMKCNS.

#### REAL I/O CONTROL BLOCKS

In order to control the activity of the I/O devices of the system and schedule I/O requests upon them, I/O control uses several types of control blocks. These blocks can be separated into two basic types:

- Static blocks that describe the components of the I/O system.
- The dynamic blocks that represent active and pending requests for I/O operations.

The I/O components of the real system are described by one control block for each channel, control unit, and device available to the control program. Units present but not represented by control blocks are not available for either user initiated or control program initiated operations.

RCHBLOK: For each channel attached to the system there exists a Real Channel Control Block (RCHBLOK) which contains:

- The channel portion of the address of its attached units,
- Status flags reflecting its availability for scheduling.
- A two-way queue anchor pointing to the list of I/O requests waiting for its services.

In addition, each RCHBLOK contains 32 half-word indexes, arranged in ascending address order, that represent the displacement into the Real Control Unit table of the control blocks for the control units attached to the channel. The 32 entries are required because the control unit address may be made up of 5 bits from the unit address. To locate the control block for a given unit, it is only necessary to:

- Index into the table in the RCHBLOK a displacement equal to twice the control unit address.
- Load the index value.
- Add the value to the base address of the Real Control Unit Table.

RCUBLOK: The Control Unit Table is composed of Real Control Unit Blocks (RCUBLOK), one for each Control unit on the system. These blocks are similar to the RCHBLOK in that they contain the control unit portion of the address and status flags, and a pointer to a queue of I/O requests. In addition the RCUBLOK contains a pointer to the RCHBLOK for the channel to which it is attached. The RCUBLOK contains a table of 16 halfword entries that represent the displacement into the Real Device Table of its attached devices. This table is referenced in the same manner as the table in the RCHBLOK.

RDEVBLOK: Each device in the system is represented by a Real Device Control Block (RDEVBLOK), contains the device portion of the unit address and status flags similar to those in RCHBLOK and RCUBLOK. There is also a pointer for those operations that are waiting for the device to become available. Fields that appear in the RDEVBLOK and not in the other blocks include a pointer to the I/O request that is currently active on the device, SIO counts, and a pointer to error and sense information. The RDEVBLOK contains a pointer to the RCUBLOK for the control unit to which it is attached and fields of device dependent information which do not affect the operation of I/O control.

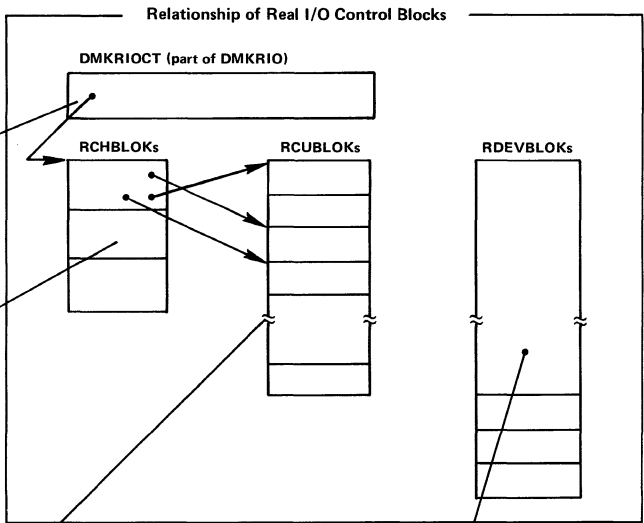
# Diag. 1B4.0 Real I/O Control Blocks

The real machine configuration is represented by a set of related control blocks. These blocks are:

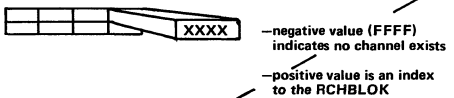
- part of the VM/370 nucleus
- built from macros in assembly of DMKRID
- loaded at system IPL and initialized then for operation.

There is one control block per channel, per control unit, and per device.

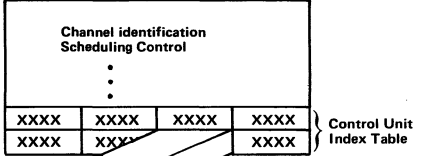
- The characteristics of VM/370 real I/O control are:
- Block multiplexing (BMPX) with RPS (Rotational Position Sensing) is used.
  - Multi-path scheduling (2 channel switching) is not used.
  - All I/O operations are handled by VM/370 scheduling and interrupt handling.



**DMKRIOCT – Real Channel Table**  
See Appendix X for a complete description

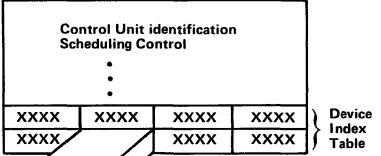


**RCHBLOK – Real Channel Block**



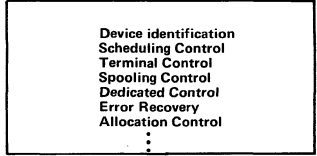
XXXX if negative (FFFF), no control unit exists  
if positive, that value is an index to the RCUBLOK

**RCUBLOK – Real Control Unit Block**



XXXX if negative (FFFF), no device exists  
if positive, that value is an index to RDEVBLOK

**RDEVBLOK – Real Device Block**



Part of the RDEVBLOK pertains to functions that are device independent; that part of the RDEVBLOK is used in the same way for all devices. However, some of the fields in the RDEVBLOK have multiple uses, depending on the device type and function.

**IOBLOK:** I/O requests that are active in the system are represented by IOBLOKs. There is one IOBLOK for each operation (that is, channel program) to be executed. The IOBLOK is constructed by the requesting task and contains such information as:

- The identity of the requestor
- The address of the channel program to be executed
- The address to which control is to be returned upon completion of the operation

In addition, the IOBLOK contains status flags that indicate the current state of the operation (such as, whether or not an error has occurred, if an Error Recovery Procedure (ERP) is in control, and the condition returned from the SIO) and the CSW associated with the interrupt that signals the end of the operation. Since IOBLOKs are queued off various I/O control blocks, they also contain forward and backward queue pointers. DMKIOS builds in them the real device address of the unit on which the operation is started.

In general, the IOBLOK representing a given operation progresses through the system by being queued, in turn, from device, control unit, and channel blocks until a path is at last free to the device. A SIO is then issued. After the operation is complete, the IOBLOK is dequeued from the RDEVELOK and stacked on a queue maintained in the dispatcher, DMKDSP. Each time the dispatcher is entered, the entries on the queue are unstacked and control is passed to the point specified in the Interrupt Return Address (IOBIRA). After I/O control stacks the IOBLOK for the given task, it attempts to restart all of the components that have been freed by the completion of the operation.

#### I/O COMPONENT STATES

The I/O components represented by the control blocks described in the section "Real I/O Control Blocks" are in one of four states and the state is indicated by the flag bits in the block status byte. If the component is not DISABLED, it is either BUSY, SCHEDULED, or AVAILABLE.

If the DISABLED bit is on, the component has been taken offline by the operator or the system and is at least temporarily unavailable. A request to use a disabled component causes the IOBLOK to be stacked with an indication of condition code 3 on the SIO and the real SIO is not performed.

A component is BUSY if it is transferring data (in the case of a channel or control unit), or if it is in physical motion (in the case of a device). If a component is BUSY, the IOBLOK for the request is queued from the control block representing that component.

A component is SCHEDULED if it is not BUSY but will become BUSY after a higher level component in the subchannel path becomes available and an operation is started. For example, if a request is made to read from a tape drive and the drive and control unit are available, but the channel is BUSY, the IOBLOK for that request is queued from the RCHBLOK for the BUSY channel and the RCUBLOK and RDEVBLOK of the drive and control unit are marked SCHEDULED. Future requests to that drive are queued from the RDEVBLOK for the SCHEDULED device. When the channel completes the operation, the next pending operation is dequeued and started; the SCHEDULED control unit and device are then marked BUSY.

The IOBLOKs for various I/O requests indicate the status of that request by a combination of the status bits in the IOBLOK and the queue in which the block resides. In general, an IOBLOK is queued from the control block of the highest level component (taken from device up to channel) in the subchannel path that is not available. Once the I/O operation is started, the IOBLOK is chained from the active IOBLOK pointer (RDEVAIOB) in the Real Device Control Block. Flags in the IOBLOK status fields may also indicate that a unit check has occurred, that a sense is in progress, or that a fatal I/O error (unrecoverable) has been recognized by ERP. After I/O control releases control of the IOBLOK, it is stacked on the queue of IOBLOKs and CPEXBLOKs anchored at DMKDSPRQ in the dispatcher and control is passed to the second level interrupt handler whose address is stored in IOBIRA.

## I/O INTERRUPTS

I/O interrupts are usually either synchronous or asynchronous. Asynchronous interrupts indicate the change in status of an I/O component from the not-ready to ready state or busy to not-busy state. In either case, if the affected component has any pending requests queued from its control block, they are restarted and whether or not the given interrupt is processed any further depends upon the status of the interrupting component. Channel available and control unit end type interrupts restart the interrupting component. An asynchronous device end is passed to the user if the device is dedicated; otherwise, the device is restarted.

An interrupt is considered to be synchronous if the interrupting device has a nonzero pointer to an active IOBLOK. In this case the processing that occurs is as follows:

- If a unit check has occurred, a SENSE is scheduled, and when the SENSE is completed, the appropriate ERP is called.
- If an ERP is currently in control of the task (indicated by a flag in the IOBLOK), return the IOBLOK to the appropriate ERP.
- If the operation is incomplete (for example, channel end is received without device end), the IOBLOK is copied and the copy is stacked but the original IOBLOK remains attached to RDEVIOB to receive the final interrupt; then, the control unit and the channel is restarted.
- If the operation is complete (that is, the device is available), the IOBLOK is unhooked from the device and stacked, and the device, control unit and channel are restored.

The restart operation usually dequeues the next IOBLOK that is queued to the restarted component and queues it to the next higher component in the subchannel path. When the channel level is reached, a SIO is issued and exit is taken to the dispatcher after handling any non-zero condition codes as previously described.

## DASD Error Recovery, ERP (DMK DAS)

Error recovery is attempted for VM/370 control program initiated I/O operations to its supported devices and for user-initiated operations to control program supported devices which use a diagnose interface. The primary control blocks used for error recovery are the RDEVBLOK, the IOBLOK and the IOERBLOK. In addition, auxiliary storage is sometimes used for recovery channel programs and sense buffers.

The initial error is first detected by the I/O interrupt handler which performs a SENSE operation if a unit check occurs. Unit check errors are then passed to an appropriate ERP. If a channel check is encountered, the channel check interrupt handler determines whether or not retry is possible and pass control to an ERP through the I/O interrupt handler. DASD errors are processed as described below.

### CHANNEL ERRORS

- Channel control check is treated as seek check. It is retried 10 times.
- Interface control check is treated as seek check. It is retried 10 times.
- Channel data check is treated as data check. It is retried 10 times.

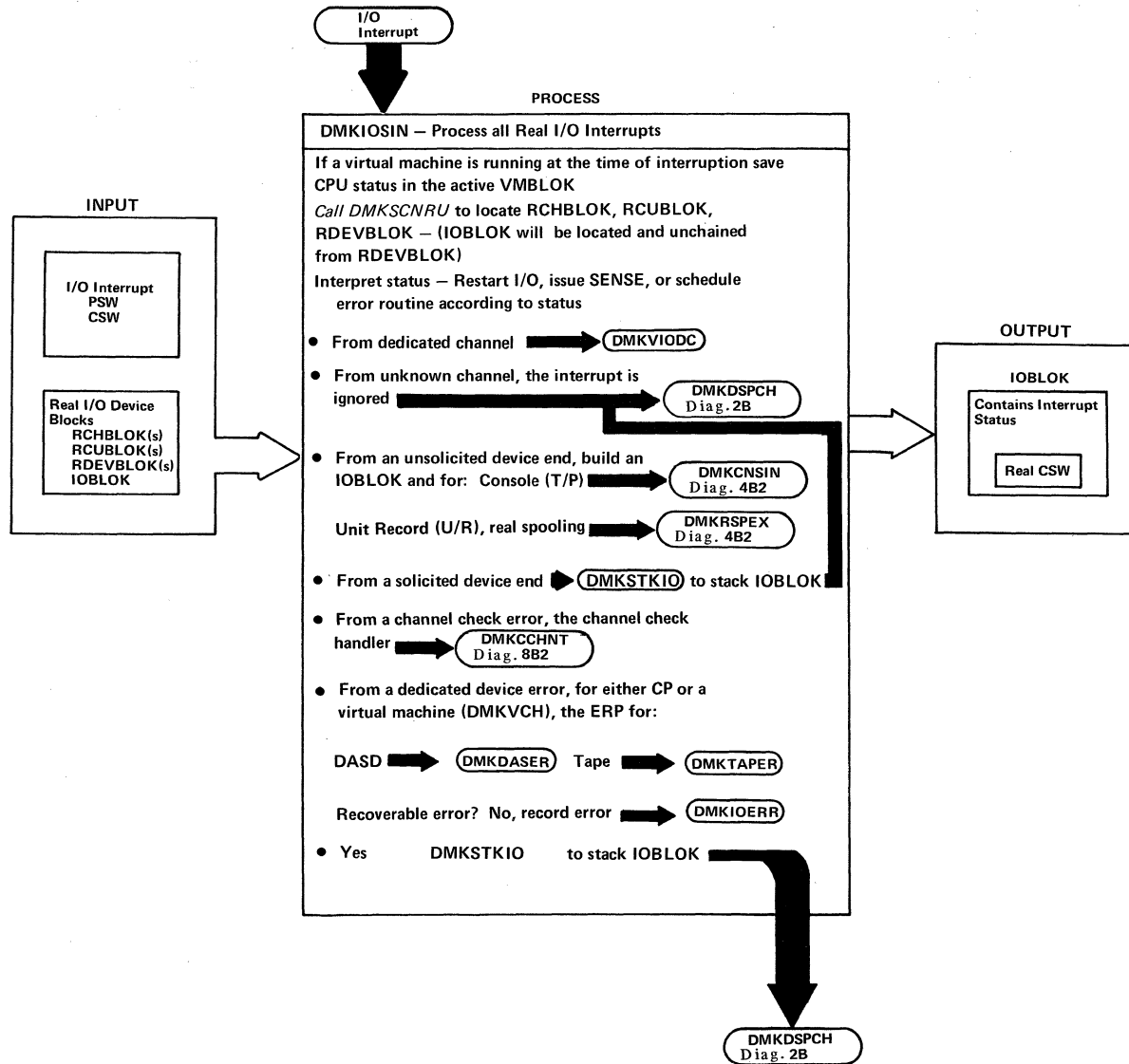
### UNIT CHECK ERRORS

Equipment check: Retry the operation once.

No record found and missing address marker: Recalibrate and retry the channel program 10 times.

No record found: Execute a READ HOME ADDRESS and check home address against seek address. If they are the same, consider the error permanent. If they are not equal recalibrate and retry the channel program 10 times.

Diag. 1B4.1. I/O Interrupt Handler





Seek check: Retry the operation 10 times.

Intervention required: Issue a message to console and wait for solicited device end. This procedure will be repeated once.

Bus out check: Retry the operation 10 times.

Data checks: Retry the operation 256 times, with a recalibrate being executed every 16th try.

Overrun: Retry the operation 10 times.

Missing address marker: Retry the operation 10 times.

Command reject: Retry the operation once.

Chaining check: Retry the operation 10 times.

Environmental data present: Issue a buffer unload command and retry the operation.

Track condition check: This error should not occur. The VM/370 control program does not use alternate tracks in its paging or spooling management. When a disk pack is formatted, any track that is marginal is marked as permanently allocated and, therefore, made unavailable for use by the VM/370 control program.

The error recovery routine keeps track of the number of retries in the IOBRCNT field of the IOBLOK. This count is used to determine if a retry limit has been exceeded for a particular error. On initial entry from DMKIOS for an error condition, the count is zero. Each time a retry is attempted the count is increased by one.

The ERP preserves the original error CSW and sense information by placing a pointer to the original IOERBLOK in the RDEVBLOK. Additional IOERBLOKS, which are received from DMKIOS on failing restart attempts are discarded. The original IOERBLOK is thus preserved for recording purposes.

If the specified number of retries fails to correct the error situation, the operator is notified and control is returned to DMKIOS. DMKIOS is notified of the permanent error by posting the IOBLOK (IOBSTAT=IOBFATAL). The error is recorded by DMKIOS via DMKIOERR.

If the error is corrected by a restart, the temporary or transient error is not recorded. Control is returned to DMKIOS with the error flag off.

Before returning control to DMKIOS on either a permanent error or a successful recovery, the ERP frees all auxiliary storage gotten for recovery CCWs, buffers, and IOERBLOKS.

The DMKIOS interface with the ERP uses the IOBSTAT and IOBFLAG fields of the IOBLOK to determine action required when the ERP returns to DMKIOS.

When retry is to be attempted the ERP turns on the restart bit of the IOBFLAG field. The ERP bit of IOBSTAT field is also turned on to indicate to DMKIOS that the ERP wants control back when the task has finished. This enables the ERP to receive control even if the retry was successful and allows the freeing of all storage gotten for CCWs and temporary buffers. The IOBRCAW is set to the recovery CCW string address.

In handling an intervention required situation, the ERP sends a message to the operator and then waits for the device end to arrive. This is accomplished by a return to DMKIOS with the ERP bit in the IOBSTAT field set on and the IOESTRT bit in the IOBFLAG field set off. When the device end interrupt arrives, the original channel program which was interrupted is then started.

The ERP flags of the IOERBLOK are also used to indicate when special recovery is being attempted. For example, a READ HOME ADDRESS command when a no record found error occurs.

The other two indications are self explanatory and are explained in Figure 3.

Field			Action to be Performed by DMKIOS
IOESTAT	IOBFLAG	IOBSTAT	
IOERP	IOBRSTRT	IOBFATAL	
1	0	0	Return control when solicited device end arrives
1	1	0	Restart using IOBRCAW
0	0	1	Permanent I/O Error
0	0	0	Retry successful

Figure 3. Summary of IOB Indicators

If the error is uncorrectable or intervention is required, the ERP calls DMKMSW for operator awareness. The specific message is identified in the MSGPARM field of the IOERBLOK.

Tape Error Recovery, ERP (DMKTAP)

Error recovery is attempted for user-initiated tape I/O operations to VM/370 control program supported devices that use the diagnose interface. The primary control blocks used for error recovery are the RDEVLOK, the IOBLOK, and the IOERBLOK. In addition, auxiliary storage is used for recovery channel programs (repositioning and erase).

The interrupt handler, DMKIOS, performs a SENSE operation when a unit check occurs. Tape errors are then passed to this DMKTAP. The sense information associated with a unit check is contained in the IOERBLOK. If a channel check is encountered, the channel check interrupt handler determines if retry is possible and passes control to the ERP through the I/O Interrupt Handler.

When an error is encountered and ERP receives control, DMKTAP determines if this the first entry into the ERP

for this task. The IOBRCNT (IOB error count) field of the IOB is zero. On this first entry, the pointer to the IOERBLOK is placed in the RDEVIOER field of the RDEVBLOK. This preserves the original error CSW and sense information for recording. Thereafter, IOERBLOKS are discarded before a retry is attempted or a permanent error is passed to IOS.

The ERP looks for two other specific conditions. If the error count field is not zero, entry must be due to a recovery attempt. Thus, it may be a solicited device end to correct an intervention required condition or a retry attempt for either tape repositioning or channel program re-execution.

The ERP keeps track of the number of retries in the IOBRCNT field of the IOBLOK to determine if a retry limit has been exceeded for a particular error. If the specified number of retries fails to correct the error, the error is recorded and DMKIOS is notified of the permanent error by turning on a status flag in the IOBLOK (IOESTAT=IOBFATAL).

If the error is corrected by DMKTAP, the temporary error is not recorded and control is returned to DMKIOS with error flags all off. When repositioning is required to attempt recovery, additional flags (EPPFLAGS) are contained in the IOERBLOK to indicate paths for specific errors (that is, data check on write must reposition, erase, and then reissue original channel program).

All error recovery is started the same except for intervention required errors. The IOBFLAG is turned on to indicate RESTART (IOBFLAG=IOBRSTRT), and the IOBRCAW (IOBLOK Restart CAW) is filled with the restart channel address word. In addition, an IOBSTAT flag is turned on to indicate that the ERP is in control so that control can be returned to ERP during all tape error recovery (IOBSTAT=IOERP). In the case of an intervention required error, the ERP sends a message to the operator, and then returns to DMKIOS with indications that tell DMKIOS the ERP is waiting for a device end on this device. This is done by clearing the restart flag and returning to DMKIOS with only the IOERP flag on.

When ERP has determined a permanent error situation or successfully recovered from an error, all auxiliary

storage gotten for recovery CCWs, buffers, and IOERBLOKs is freed before a return is made to DMKIOS (see Figure 3 for a summary of the IOB indicators).

If the error is uncorrectable or operator intervention is necessary, the ERP calls the message writer to write the specific message.

### Virtual I/O Interrupts

When an I/O interrupt is received (see Figure 4), the IOBLOK is stacked for dispatching and control is passed to the address specified in the IOBIRA (Interrupt Return Address) field. For operations requested by DMKVIOEX, the return address is DMKVIONT (Virtual Interrupt Return Address). When DMKVIONT receives control from the dispatcher, it loads the virtual address of the unit with which the interrupt is associated from the IOBLOK and calls DMKSCNVU to locate the virtual device control blocks. DMKVIONT then tests the IOBLOK status field to determine the cause for the interrupt. If the block has been unstacked due to an interrupt, the field is zero. If the operation was not started, it contains the condition code from the real SIO.

Note: The VIRA should not see a real condition code 2 as the result of a SIO, since channel busy conditions are detected and reflected before any real I/O operation is attempted.

A condition code 3 is reflected to the user and exit is taken to the dispatcher. For condition code 1, the CSW status field in the IOBLOK is examined to determine the cause for the CSW stored condition. The status is reflected to the user and various components of the virtual configuration may be freed, if the status so indicates. For example, if the CSW status indicated both channel end and device end, the operation was immediate and has completed. Thus, the CCW string (real) may be released and all virtual components marked available.

The CSW status returned for a virtual interrupt must be tested in the same manner, with the additional requirement that the status be saved in the affected

virtual I/O control blocks and that the CSW be saved in the VDEVCSW field for the device causing the interrupt. If the unit check bit is on in the status field, the sense information saved in the associated IOERBLOK (pointed to by the IOBLOK) must be retained so that a sense initiated by the virtual machine receives the proper information.

In any case, when an interrupt is received for a virtual device, a bit in the interrupt mask, VCUDVINT, for the device's control unit is set to one. The bit that is set is the one corresponding to the relative address of the interrupting device on the control unit. For example, if device 235 interrupts, the fifth bit in the VCUDVINT mask in the VCUBLOK for control unit 30 on channel 2 is flagged. Similarly, the bit in the VCHCUINT in the affected VCHBLOK is also set; in this case, bit 3 in VCHBLOK for channel 2. If the interrupt is a channel class interrupt (PCI or CE), the address of the interrupting unit (235) is stored in the VCHCEDEV field in the VCHBLOK. The final interrupt flag is set in the VMPEND field in the VMBLOK for the interrupted user; the bit set corresponds to the address of the interrupting channel. The next time, the user is dispatched and becomes enabled for I/O.

### Scheduling I/O Requests

A task that requests an I/O operation must specify the device on which the operation is to take place and must provide an IOBLOK that describes the operation. Upon entry to DMKIOS, Register 10 must point to the IOBLOK. The IOBLOK must contain at least a pointer to the channel program to be started in IOBCAW and the address to which the dispatcher is to pass control in IOBIRA. In addition, the flags and status fields should be set to zero. If the operation is a VM/370 control program function such as spooling or paging, the entry point DMKIOSQR is called. If the requestor is the virtual I/O executive (OMKVIDEX) attempting to start a user operation, the entry point DMKIOSQV is called and some additional housekeeping is done. In either case, an attempt is made to find an available subchannel path from the device to its control unit and channel. If a component in the path is BUSY or SCHEDULED, the IOBLOK for the request is queued to the control block of the component.

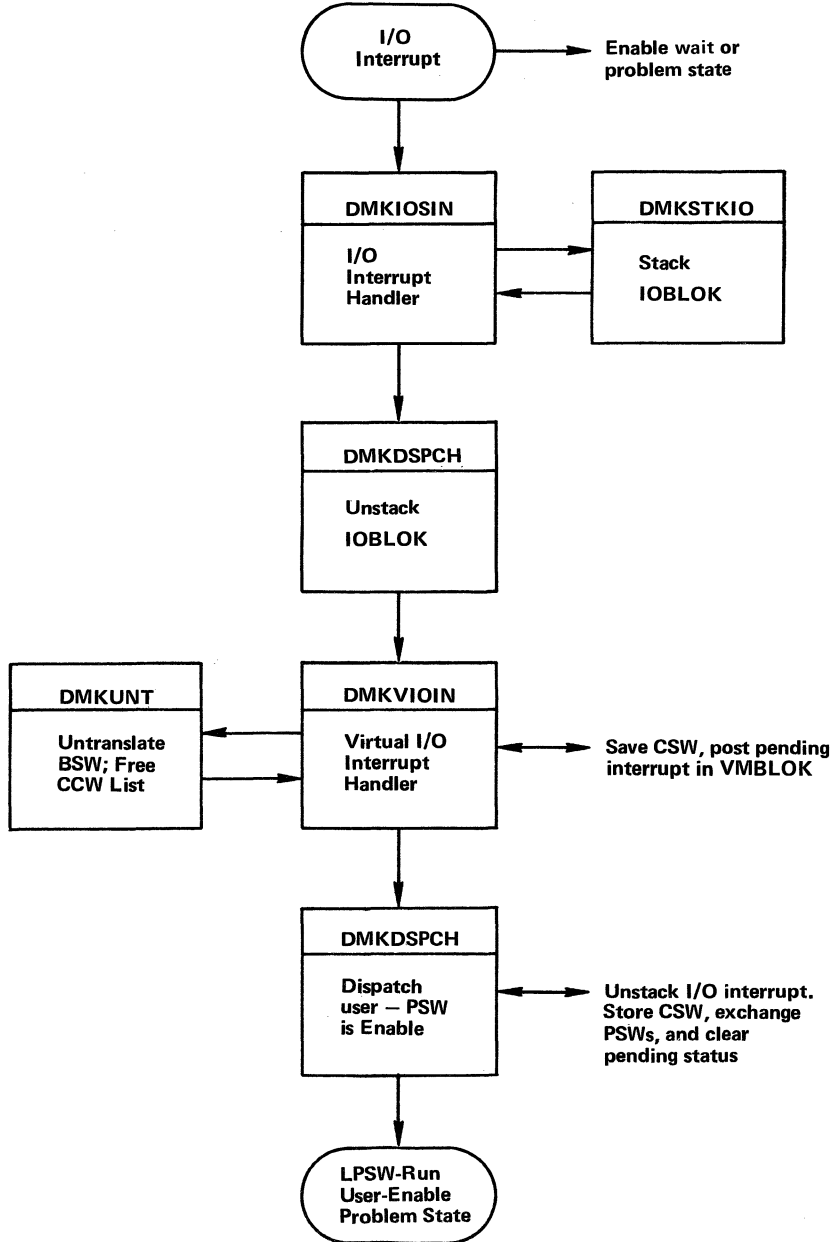
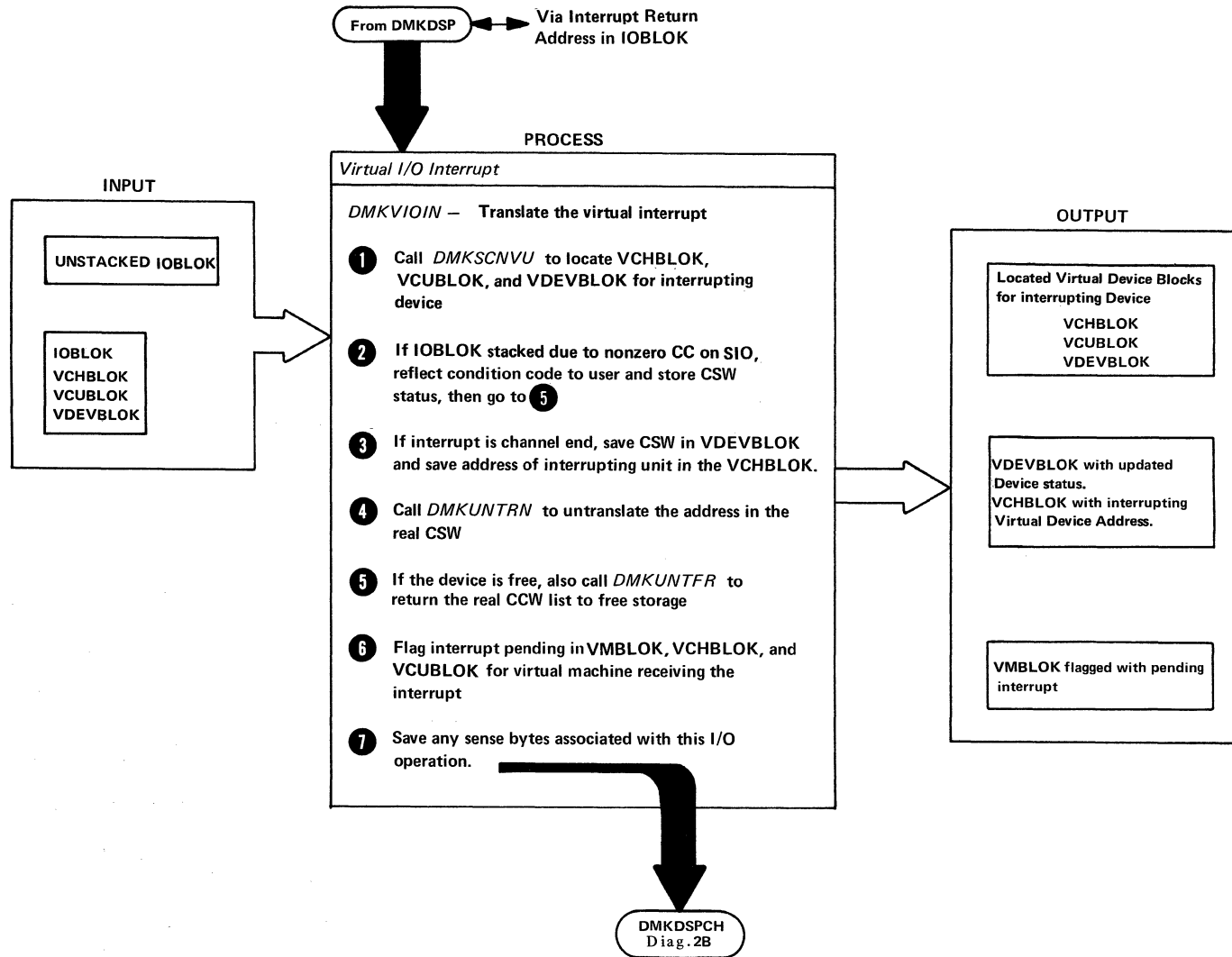


Figure 4. Overview of a Virtual I/O Interrupt

Diag. 1B4.2. Virtual I/O Interrupt



Requests are usually queued first-in first out (FIFO), except those:

- To moveable head DASD devices that are queued in order of seek address.
- That release the affected component after initiation (SEEKS and other control commands) which are queued last-in first out (LIFO) from the control block.

Regardless of whether or not the operation has been successfully started, the caller requesting the I/O operation receives control back from DMKIOS. If a free path to the device is found, the unit address is constructed and an SIO is issued. If the resulting condition code is zero, control is returned to the caller; otherwise, the code is stored in the requestor's IOBLOK along with any pertinent CSW status, the IOBLCK is stacked, any components that become available are restarted, and control is returned to the caller.

Ordered Seek Queueing: Requests to start I/O on system devices are normally handled FIFO. However, requests to moveable head DASD devices are queued on the device in ascending order by seek address. This ordered seek queueing is performed to minimize intercylinder seek times and to improve the overall throughput of the I/O system.

The VM/370 control program assumes that very few virtual machines will do chained seeks; hence, the first logical address represents where the arm will be positioned upon completion of the I/O operation. Ordered seek queueing is based on the relocated real cylinder. DMKIOS uses the cylinder location supplied in IOBCYL for ordered seek queueing. This field is initialized by the calling VM/370 control program routine for paging and spooling or by the CCW translator for virtual I/O. The CCW translator DMKCCW supplies the IOBCYL value in the following manner.

- Read IPL record, relocated to virtual cylinder 0
- Recalibrate, issue a real recalibrate and then seek to virtual cylinder 0
- Channel seeks, relocate to the virtual cylinder

The IOBLOK queueing subroutine of DMKIOS recognizes

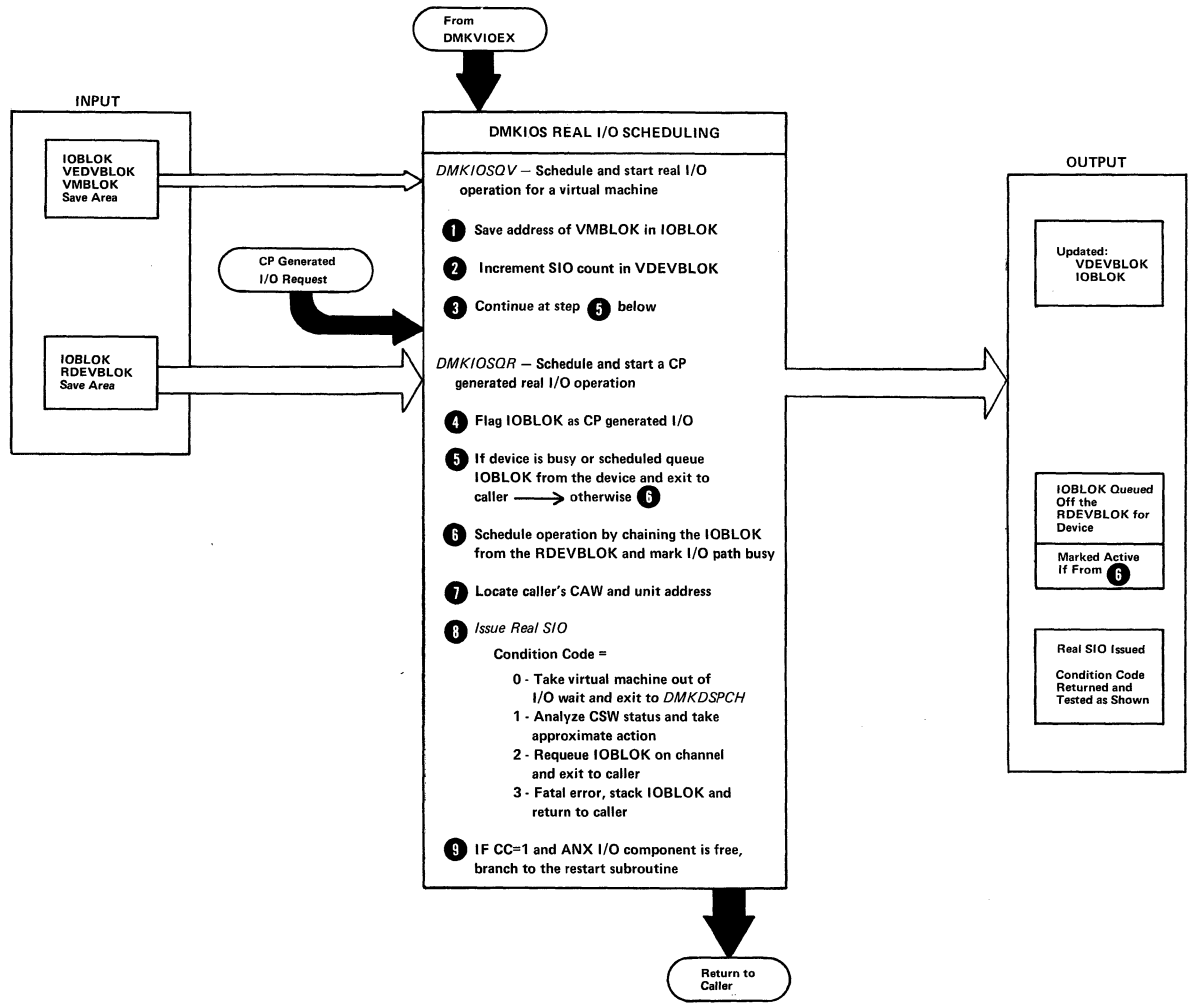
that a request is being queued on a moveable head DASD device by means of the device class and type fields of RDEVBLK. Instead of adding the IOBLOK to the end of the queue on the RDEVBLK, the queueing routine sorts the block into the queue based on the cylinder number for the request. The cylinder number for any request to a DASD device is recorded in the field IOBCYL. The queue of IOBLOKs on a real device block is sorted in ascending order by seek address, unless the entire device is dedicated to a given user. In this case, DMKIOS does not automatically schedule the device, and no more than one request can be outstanding at any one time.

When an outstanding I/O request for a device has completed, DMKIOS attempts to restart the device by dequeuing and starting the next IOBLOK queued on the device. For non-DASD devices, this is the first IOBLOK queued. However, for moveable head DASD devices, the queued requests are dequeued in either ascending or descending order, depending on the current position (recorded in RDEVCYL) and the direction of motion of the arm. If the arm is seeking up (that is, toward the higher cylinder numbers), the queue of IOBLOKs is scanned from the first block toward the last until an IOBLOK is found with an IOBCYL value equal to or greater than the value in RDEVCYL, or until the end of the queue is reached. At this point, the device is flagged as seeking down and the queue is scanned from last to first until an IOBLOK with an IOBCYL value equal to or less than RDEVCYL is found. When IOBLOK is found, it is dequeued and started. The direction of motion is remembered in an RDEVFLAG bit and the next request is dequeued in the down direction until the head of the queue is reached.

Because the queue itself is a two-way chained list, no special handling for null or unity set lists is required, and the ordered seek algorithm returns to FIFO queueing.

Dedicated Channel Support: One of the facilities of the VM/370 control program allows a virtual machine to control one or more channels on a dedicated basis. The channels are attached to the virtual machine by using the privileged ATTACH CHANNEL command. A virtual machine can have one or more dedicated channels. In addition, channels can be split between virtual machines but a dedicated channel cannot be shared between two virtual machines. For instance, channel 1

Diag. 1B4.3. I/O Scheduling



could be dedicated to virtual machine A, and channel 2 could be dedicated to virtual machine B, or they could be both dedicated to virtual machine A or B.

With a dedicated channel, all virtual machine device addresses must be identical to the real machine device addresses. For instance, virtual device 130 must be on real device 130, and virtual device 132 must be on real device 132. With dedicated channels, the VM/370 control program does not perform any virtual device address mapping. With a dedicated channel in effect, a virtual machine I/O operation to one of the dedicated devices on that channel results in the control program performing the operation directly on that device and reflecting the true condition code back to the virtual machine. None of the I/O operations are passed through control program's normal channel scheduling since the channel is completely dedicated to the virtual machine and any conditions in the channel are a direct result of that virtual machine's operation of that channel.

It is expected than any I/O new PSW for a virtual machine operating system has all channels masked off. Thus, when the VM/370 control program receives a hardware interrupt from a dedicated channel it immediately disables all further interrupts on that channel. The interrupt is then reflected to the virtual machine. The real channel stays disabled until the virtual machine issues an instruction to enable that channel. At that time, the VM/370 control program performs a hardware function to enable the real channel.

By using the dedicated channel feature, a virtual machine bypasses the VM/370 control program overhead associated with channel scheduling and virtual machine interrupt stacking. The channel scheduling is bypassed by performing the I/O operation directly and the interrupt stacking is bypassed by disabling the channel and having the hardware perform the true interrupt stacking.

The VM/370 control program error recording and channel recovery procedures are still in effect for dedicated channels. The dedicated channel support can be used in conjunction with the Virtual=Real feature for any virtual machine that is occupying the Virtual=Real storage space.

#### DISPATCHER/SCHEDULER

The module that selects dispatchable users from the population is DMKSCH, the Scheduler. The module that tests and alters the resources of the CPU is DMKDSP, the Dispatcher. The auxiliary routine that assists the Scheduler and Dispatcher is the request stack maintenance routine, DMKSTK.

In order to make decisions on both dispatching and scheduling, the control program classes all users into various categories, and recognizes user machines as being in one of several states. The user categories recognized are classed as being either interactive or non-interactive.

- An interactive user is one whose use of the system is punctuated by regular and frequent terminal I/O, and does not execute long CPU loops. A user becomes eligible to enter interactive status whenever a channel program for virtual console I/O has completed, or whenever I/O for a dedicated or dialed virtual telecommunications line has completed.
- A non-interactive user is one who has violated an interactive criterion, or one who has entered an idle wait state by entering console function mode (equivalent to stopped state), or by loading a wait state PSW that is not enabled for any busy channel. The control program schedules interactive users ahead of non-interactive users. Non-interactive users are subdivided into several classes. Normal non-interactive users are scheduled via a priority scheme described below. A user is allowed to execute for a specified time period and he is then placed in a list of those users who are waiting.

In order to give preference to certain classes of users, a priority scheduling scheme allows users to be scheduled with a priority class. The priority is a number assigned by the directory; however, the number may be altered by the system operator.



## USER DISPATCHING LISTS AND MACHINE STATES

In order to efficiently manage the large inventory of potential users that are logged on to the system, the control program defines several states that a virtual machine may occupy. The scheduler can move a virtual machine from one state to another; however, a virtual machine may exist in only one state at any given instant. The control program can then make scheduling and dispatching decisions by looking only at the subset of users that are in the appropriate state. To facilitate this search, it also maintains lists of users in certain executable states.

A user's virtual machine may be in one of the following states:

<u>State</u>	<u>Meaning</u>
1	Interactive and dispatchable (in queue1, in DISPATCH list)
2	Interactive and not dispatchable (in queue1, not in DISPATCH list)
3	Interactive and eligible for queue1, but queue1 is full (waiting for queue1, in ELIGIBLE list)
4	In wait state with terminal read or write active
5	Non-interactive and dispatchable (in queue2, in DISPATCH list)
6	Non-interactive and not dispatchable (in queue2, not in DISPATCH list)
7	Non-interactive and eligible for queue2, but queue2 is full (waiting for queue2, in ELIGIBLE list)
8	Idle - waiting for asynchronous I/O or external interrupt, or stopped (in Console Function Mode)

Two lists of users are maintained by the scheduler:

- The DISPATCH list
- The ELIGIBLE list

Entries on the DISPATCH list are the VMBLOKS for those users in states 1 and 5, and represent the users that can be run at any given time. The DISPATCH list is sorted by dispatching priority, which is the ratio of CPU time to wait time over the life of the current user task. A task is defined as that execution which takes place between terminal reads or entry to enabled wait (that is, movement from state 4 or 8 to state 1) and is re-projected for a user each time he is dropped from a queue. Users entering state 1 always have a priority of 0.

The ELIGIBLE list is composed of those users in states 3 and 7; these users are potentially executable but due to the current load on the system they are not allowed to compete for the CPU. As soon as a user in the DISPATCH list is dropped from queue, the highest priority user(s) in the ELIGIBLE list is added to the DISPATCH list, subject to the restriction that his projected working set must not exceed the remaining system capacity. The ELIGIBLE list has two components; a section composed of those virtual machines waiting for Q1 (interactive) and a section composed of those virtual machines waiting for Q2 (non-interactive). Each section of the list is sorted by scheduling priority, which is determined at the time the virtual machine is added to the ELIGIBLE list, as follows:

1. The virtual machine's projected working set size, calculated the last time it was dropped from a queue, is expressed as a percentage of the amount of main storage available for paging. This percentage, usually between 0 and 100, is multiplied by the Paging Bias Factor (stored at DMKSCHPB).
2. The virtual machine's user priority (the priority set by the directory or the class A "SET PRIORITY" command) is multiplied by the User Bias Factor (stored at DMKSCHUB), and is added to the Paging Bias calculated in step 1.
3. The sum of Paging and User Bias is divided by the sum of the Bias Factors to obtain a weighted average.
4. A base priority is obtained by storing the TOD clock and using the high order word, which increments by 1 approximately once per second. This word is then modified by shifting it left or

right based on the Priority Delay Factor (stored at DMKSCHPD). If DMKSCHPD is positive, it indicates a right shift, thereby increasing the delay interval of the base priority; while a negative value indicates a left shift.

5. The weighted average obtained in step 3 is then logically added to the adjusted base obtained in step 4.
6. If the virtual machine is entering Q2 for the first time after being dropped from Q1, the Interactive Bias Factor (stored at DMKSCHIB) is subtracted from the priority obtained in step 5. If the virtual machine is entering Q1, or if it was last dropped from Q2, the Interactive Bias is not applied.
7. The result of steps 1 through 6 is the scheduling or eligible list priority, and is stored in the VMEPRIOR field of the VMBLOK.

The VMBLOK is then sorted into the appropriate section of the ELIGIBLE list in ascending value of VMEPRIOR. The effects of the various biases and the delay factor are illustrated by the following examples.

1. Assume that two virtual machines are to be added to the ELIGIBLE list for Q2. The Paging Bias Factor is 1, the User Bias Factor is 1, and the Priority Delay Factor is 0. Virtual machine "A" has a projected working set size of 80 percent of available storage and a user priority of 50. Virtual machine "B" has a projected working set size of 20 percent of available storage and also has a user priority of 50. The biases are obtained as follows:

User	Paging Bias	User Bias	Weighted Bias
A	80 X 1	+ 50 X 1	= 130/2 = 65
B	20 X 1	+ 50 X 1	= 70/2 = 35

If "A" is added to the eligible list at base time 0, its eligible list priority will be 65. If the Priority Delay Factor is 0, "B" will be added ahead of "A" provided that "B" is eligible for entry to the list within the next (65-35) 30

seconds. If the Priority Delay Factor is set to +1, the base will be incremented once every two seconds. Therefore, although the bias difference is still 30, the delay time is now 60 seconds.

2. In order to force "A" to be given a weighted bias equal to "B," a priority differential is calculated as follows:

$$\frac{80 + A}{2} = \frac{20 + B}{2} ; A = B - 60$$

Therefore, for the biases to be equal, "A" must have a priority of 60 less than "B." For example, if "A" is given a priority of 10 and "B" is given a priority of 70, the biases would compute as follows:

User	Paging Bias	User Bias	Weighted Bias
A	80 X 1	+ 10 X 1	= 90/2 = 45
B	20 X 1	+ 70 X 1	= 90/2 = 45

3. The large difference in priorities could be lessened by increasing the User Bias Factor. If the User Bias Factor is set to 3 instead of 1, the calculated priority differential is as follows:

$$\frac{80 + 3A}{4} = \frac{20 + 3B}{4} ; 3(B - A) = 60 ; A = B - 20$$

Now, "A" requires a priority of only 20 less than "B" to achieve parity. For example:

User	Paging Bias	User Bias	Weighted Bias
A	80 X 1	+ 30 X 3	= 170/4 = 42
B	20 X 1	+ 50 X 3	= 170/4 = 42

The above examples illustrate the following general points about the use of the bias factors, the delay factor, and the user priority value:

1. The Paging and User Bias Factors are a measure of the relative importance of the bias value. A high

user bias will allow greater discrimination via the assigned priority; while a high paging bias makes storage requirement the primary scheduling parameter.

2. The user priority value, in the directory, is the means by which the paging priority may be overridden, and the means through which selected users will obtain improved performance.
3. The Priority Delay Factor is the measure of the impact which the paging and user biases are to have. The greater the delay value, the greater is the maximum delay that can be experienced by a given user.
4. The Interactive Bias Factor is a tool that enhances command response to conversational commands which require disk I/O, and which may be partially executed in Q2.

If the Paging Bias Factor is non-zero, the net effect of the priority scheme is to discriminate against users who require large amounts of real storage. This discrimination results in a higher level of multiprogramming and increased CPU utilization; however, it must be traded off against poorer throughput for large users. The distributed Scheduler is not biased; the bias factors are as follows:

```
Paging Bias Factor      (DMKSCHPB) = 0
User Bias Factor        (DMKSCHUB) = 1
Priority Delay Factor    (DMKSCHPD) = 0
Interactive Bias Factor  (DMKSCHIB) = 0
```

Thus, the basic VM/370 Scheduler will schedule virtual machines FIFO within user priority; the same algorithm provided with the basic Release 1.0 system.

Figure 5 is a graphic breakdown of the user states, showing the relationship between interactive and non-interactive states, in-queue and not-in-queue states, and in-list and not-in-list states.

	In-Queue		Not-in-Queue	
	DISPATCH List	No List	ELIGIBLE List	No List
Interactive	1	2	3	4
Non-Inter.	5	6	7	8

Figure 5. User Dispatching States

Figure 6 shows the possible user-state changes and the reasons for them; any changes not described are not possible.

#### CONTROLLING THE DEPTH OF MULTIPROGRAMMING

In order to control the number of users allowed in queue, the scheduler monitors the paging activity of all users and of the system as a whole. A decision as to whether or not to move a potential user from the eligible to the dispatch list is based upon whether or not that user's projected working set will exceed the system's remaining capacity. Individual user's working sets are calculated and projected at queue drop time according to one of the following formulas:

$$P = (A + P) / 2$$

$$\text{If } (LP - LA) * (P - A) < 0$$

-- or --

$$P = A$$

$$\text{If } (LP - LA) * (P - A) \geq 0$$

Note: See the Key for the meaning of the symbols.

The working set is added to the current system load, which consists of the sum of the working sets for all users currently in a queue. The sum is compared to the system maximum, which is equal to the number of dynamically assignable pages in the system. If the user's projected working set will not push the system load over the maximum, he is placed in the queue and added to the dispatchable list.

Status Change		Reason for Status Change
From	To	
1	2	PAGEWAIT, SIO-WAIT, or enabled wait for any busy channel
1	4	Enabled wait for interactive terminal read or write
1	5	Exceeds in-queue time slice
1	7	Same as 1 to 5 except that queue2 is full
1	8	Wait without active I/O, disabled wait, or hit ATTN
2	1	Wait condition complete
2	5,7	Wait completes, but in-queue time slice exceeded
3	1	Another user drops from queue1 and now there is room
4	1	Terminal I/O completes while user is waiting
4	3	Terminal I/O completes, but queue1 is full
5	1	Terminal I/O completes while user is active in queue2
5	4	User puts up terminal read or write and enters wait
5	6	PAGEWAIT, SIO-WAIT, or enabled wait for busy channel
5	7	Dropped from queue2 due to in-queue time-slice end
5	8	Wait without active I/O, disabled wait, or hit ATTN
6	5	Wait condition completes
7	5	Room is found in queue2
8	5,7	Asynchronous I/O or External Interrupt, or BEGIN

Figure 6. User Status Changes

**Key:**

A = Actual working set at queue drop time

- LA = Last actual working set
- LP = Last projected working set
- P = Current projected working set

The actual working set, A, is determined at queue drop time by the following formula:

$$A = \left[ \left( \sum_{i=1}^N PR_i \right) / N + \text{Steals} \right] \text{ whichever is greater}$$

— or —

$$\left[ \text{Pages referenced} \right]$$

where:

- N = Number of page reads while in queue.
- PR<sub>i</sub> = Number of pages resident at the i<sup>th</sup> page read.
- Steals = Number of times page wait was entered due to a stolen page.

The number of referenced pages is determined by scanning the user's page tables for software referenced bits. These bits are set by DMKPTRAN when the page is taken from the user by the control program. Thus the actual working set is generally the average number of pages resident at each page read. However, this estimate is sensitive to the overall system paging activity for the following reasons:

1. If there is no paging load on the system, there will be one page read for each resident page, and no steals; the working set will therefore tend to be equal to about one half of the resident page total.
2. As paging activity increases, and the working set locality shifts, the working set will tend to increase toward the average number of resident pages.
3. If paging activity becomes excessive, the number of page steals will increase to the extent that the working set will expand to the maximum of the total number of pages referenced while in the queue.

In summary, the scheduler selects the subset of logged-on users that are allowed to compete for the resources of the CPU, with the constraint that a new user is not added to the active subset if his projected main storage requirement, added to that of the other active users, causes the current capacity of the system to be exceeded. Selection within scheduling priority simply means that a executable user of high priority is always added to the active subset (to a queue) before a executable user of lower priority. If the paging bias mechanism is activated by setting the Paging Bias Factor to a non-zero value, scheduler selection will be in favor of smaller users; otherwise, selection is

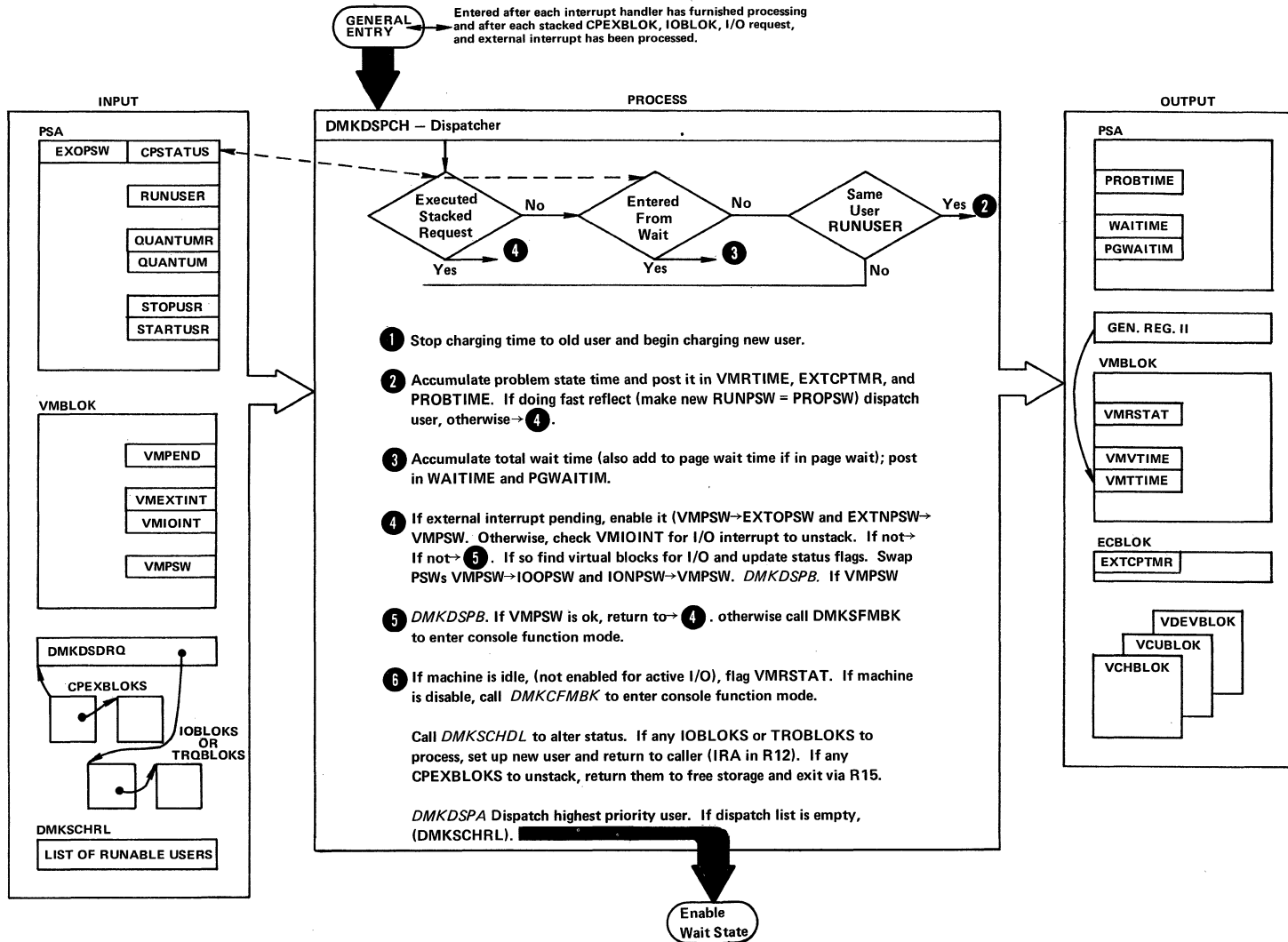
round robin within priority. Once the active subset (the set of in-queue users) has been selected, the dispatcher allocates resources of the CPU among them.

The list of executable users in a queue is sorted by dispatching (as opposed to scheduling) priority. The dispatching priority is a running average of a given user's CPU time/wait-time ratio. Thus, users who are most likely to go into wait state, based on past performance, are dispatched ahead of those whose demands on the CPU are more extensive. This simple ratio priority is normally altered if a user is identified as compute bound by means of the fact that he has executed for at least 50 ms. without entering the wait state. In this case, he is placed at the bottom of the dispatchable list. On the other hand, users identified as interactive by virtue of the frequency their requests for terminal I/O are placed at the top of the dispatchable list.

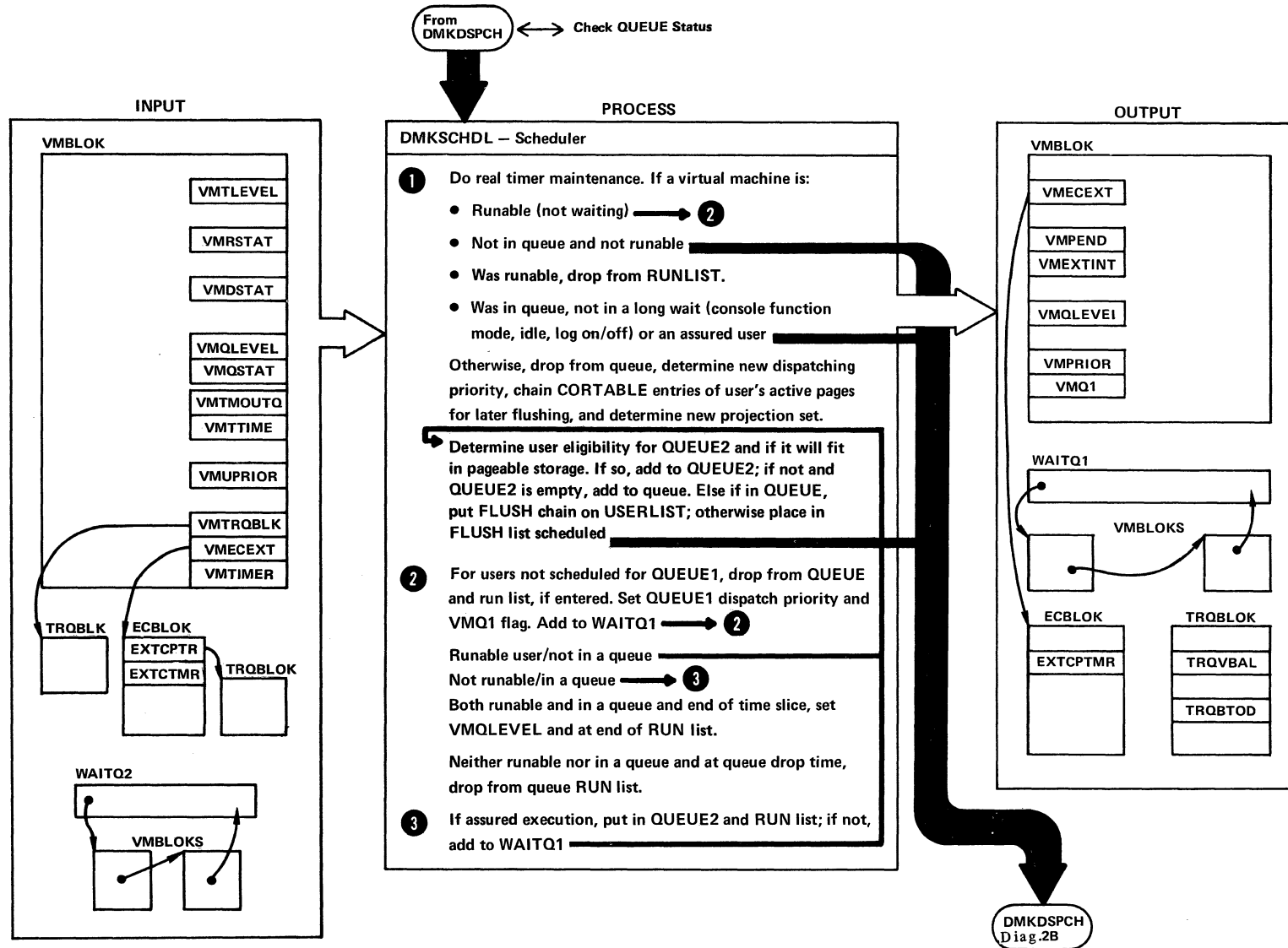
#### FAVORED EXECUTION OPTIONS

When the resources of the CPU (and real storage) are being allocated, the dispatching and scheduling functions are implemented in such a manner that options exist that allow an installation to designate certain users (virtual machines) that are to receive preferential treatment.

Diag. 2B. Dispatcher



Diag. 3B. Scheduler



The favored execution options allow an installation to modify the algorithms described above and force the system to devote more of its resources to a given user than would ordinarily be the case. The options provided are:

1. The basic favored execution option.
2. The favored execution percentage.

The basic favored execution option means that the user so designated is never to be dropped from the active (in queue) subset by the scheduler. When the user is executable, he is to be placed in the dispatchable list at his normal priority position. However, any active user represents either an explicit or implicit commitment of main storage. An explicit storage commitment can be specified by either the Virtual=Real option or the reserved page option. An implicit commitment exists if neither of these options are specified, and the scheduler recomputes the user's projected work-set at what it would normally have been at queue-drop time. Multiple users can have the basic favored execution option set. However, if their combined main storage requirements exceed the system's capacity, performance can suffer due to thrashing.

The basic favored execution option removes the primary source of elapsed time stretch-out in a loaded time-sharing environment. However, if the favored task is highly compute bound and must compete for the CPU with many other tasks of the same type, an installation can define the CPU allocation to be made. In this case, the favored execution percentage option can be selected for one virtual machine. This option specifies that the selected user, in addition to remaining in queue, receives a given minimum percentage of the total CPU time, if he can use it. The percentage is assured in the following manner:

1. The in-queue time slice is multiplied by the requested percentage and added to the user's current total CPU time usage.
2. The favored user, when he is executable, is then always placed at the top of the dispatchable list until he has obtained his guarantee.

3. If the user obtains his guarantee before the interval has elapsed, he is placed in the dispatchable list according to his calculated dispatching priority.
4. In any case, at the end of the in-queue time slice the percentage is recomputed and the process repeated.

These options can impact the response time of interactive users and only one favored percentage user is allowed at any given time.

#### DISPATCHING/SCHEDULING SUPPORT ROUTINES

Most of the routines in the CP nucleus are reentrant and multiple control program or user tasks can make use of one routine at the same time. However, there are certain areas where requests for a resource must be serialized (as in paging) or delayed while previous requests are serviced (as in requests to schedule I/O).

#### The CP Request Stack

The routine handling the request obtains a CPEXBLOK from free storage and stores the caller's registers in it; when the requested resource is free, the CPEXBLOK is stacked for the dispatcher via a call to the Request Stack Manager (DMKSTKCP). The dispatcher unstacks the block and exits to the requesting routine the next time it is entered. I/O requests are stacked in the same manner, except that the stacking vehicle is the IOBLOK, and return is passed to the address specified in the interrupt return address (IOBIRA). In either case, it should be noted that the dispatcher always unstacks and gives control to any stacked IOBLOKs and CPEXBLOKs prior to dispatching a user. This guarantees that control program information needed by a user (such as page availability) is always as up-to-date as possible.



### CP INTERNAL TRACE TABLE

CP provides an internal trace table where events that occur in the real machine may be recorded. The size of the trace table depends on the amount of real storage available at IPL time. For each 256K bytes (or part thereof) of real storage available at IPL time, one page (4096 bytes) is allocated to the CP trace table. The storage thus allocated is contiguous and each entry is 16 bytes long. The first byte of each trace table entry, the identification code, identifies the event being recorded. Events that are traced are:

- External interrupts
- SVC interrupts
- Program interrupts
- I/O interrupts
- Free storage requests
- Release of free storage
- Entry into dispatch
- Queue drop
- Run user requests
- Start I/O
- Unstack I/O interrupts
- Storing a virtual CSW
- Test I/O
- Halt device

The main initialization routine, DMKCP1, allocates storage to the CP trace table and activates internal tracing. If you do not wish to record events in the trace table, the class A or E command MONITOR STOP can

be issued to suppress recording. The pages allocated to the trace table are not released and recording can be restarted at any time by issuing the MONITOR START command. If the VM/370 system should abnormally terminate and automatically restart, the tracing of events on the real machine will be active. After a VM/370 IPL (manual or automatic), CP internal tracing is always active.

The first event traced is placed in the lowest trace table address. Each subsequent event is recorded in the next available trace table entry. Once the trace table is full, events are recorded at the lowest address (overlying the data previously recorded there). Tracing continues with each new entry replacing an entry from a previous cycle.

The trace table can be used to determine the events that preceded a CP system failure. An ABEND dump contains the CP internal trace table along with the pointers to it. The address of the start of the trace table, TRACSTRT, is at location X'0C'. The address of the byte following the end of the trace table, TRACEND, is at location X'10'. The address of the next available trace table entry, TRACCURR, is at location X'14'. The trace table entry for the last event completed is obtained by subtracting 16 bytes (X'10') from the address stored in TRACCURR.

There are fourteen possible types of trace table entries, each uniquely identified by the value of the first byte. Figure 7 describes the format of each type and identifies the CP module that records the event.

Type of Event	Module	Identification Code (hexadecimal)	Format of Trace Table Entry															
External interrupt	DMKPSA	01	X'01' 0	1	X'000000000'	6	Interrupt Code	8	External Old PSW	15								
SVC interrupt	DMKPSA	02	X'02' 0	1	GR 15	4	Instruction Length Code	6	Interrupt Code	8	SVC Old PSW	15						
Program interrupt	DMKPRG	03	X'03' 0	1	X'000000'	4	Instruction Length Code	6	Interrupt Code	8	Program Old PSW	15						
Not used		04																
I/O interrupt	DMKIOS	05	X'05' 0	1	X'00'	2	Device Address	4	I/O Old PSW + 4	8	CSW	15						
Free Storage (FREE)	DMKFRE	06	X'06' 0	1	GR 11 at entry	4	GR 0 at entry	8	GR 1 at exit	12	GR 14	15						
Return storage (FRET)	DMKFRE	07	X'07' 0	1	GR 11 at entry	4	GR 0 at entry	8	GR 1 at entry	12	GR 14	15						
Enter scheduler	DMKSCH	08	X'08' 0	1	Contents of VMRSTAT, VMDSTAT, and VMOSTAT		4	Address of VMBLOK	8	Value of CPU Timer			15					
Queue drop	DMKSCH	09	X'09' 0	1	Address of VMBLOK	4	Old Priority	6	New Priority	8	Number of Resident Pages	10	Projected Working Set	12	Number of Referenced Pages	14	Current Page load (PSA)	15
Run user	DMKDSP	0A	X'0A' 0	1	X'000000'	4	RUNUSER value from PSA	8	RUNPSW value from PSA				15					
Start I/O	DMKIOS	0B	X'0B' 0	1	Condition Code	2	Device Address	4	Address of IOBLOK	8	CAW	12	For CC = 1, CSW + 4 otherwise this field is not used			15		
Unstack I/O interrupt	DMKDSP	0C	X'0C' 0	1	X'00'	2	Virtual Device Address	4	Address of VMBLOK	8	Virtual CSW				15			
Virtual CSW store	DMKVIO	0D	X'0D' 0	1	Instruction Operation Code	2	Virtual Device Address	4	Address of VMBLOK	8	Virtual CSW				15			
Test I/O	DMKIOS	0E	X'0E' 0	1	Condition Code	2	Device Address	4	Address of IOBLOK	8	CAW	12	For CC = 1, CSW + 4 otherwise this field is not used			15		
Halt Device	DMKIOS	0F	X'0F' 0	1	Condition Code	2	Device Address	4	Address of IOBLOK	8	CAW	12	For CC = 1, CSW + 4 otherwise this field is not used			15		

Figure 7. CP Trace Table Entries

## SPOOLING

The spooling support in the VM/370 control program performs three functions.

First, to simulate the operation of the virtual unit record devices that are attached to each user's virtual machine configuration. The simulation is done in such a way that it appears to the program in the virtual machine that it is controlling a real unit record device. This support involves the interception and interpretation of user SIOs, the movement of data to and from the user's virtual storage space, and the reflection of the necessary interrupt codes and ending conditions in PSW's, CSW's and sense bytes. This support is provided by the Virtual Spooling Executive.

Second, to operate the real unit record equipment attached the system that is used to transcribe user output spool files from input from the real card readers. This function is provided by the Real Spooling Executive.

Third, to provide an interface among the users, the system operator, and the spooling system so that the location, format, priority and utilization of the systems spooling data and resources can be controlled.

### SPOOL DATA FORMAT

#### Spool Buffer Format

The buffers used for collecting and writing spool data are all one page (4096 bytes) in length, and contain both the data to be transcribed and all CCWs necessary for operating the unit record devices that perform the transcription. The data is provided in the exact format required with no compression except that trailing blanks are suppressed. The first two doublewords of each buffer contain linkage information described below, followed by the data and CCWs.

Each spool logical record (card or print line) is stored as one data moving CCW (READ or WRITE), a TIC to

the following CCW, and the full data record. Space is left at the end of each buffer so that a SENSE command can be inserted in order to force concurrent channel end and device end. For card punch channel programs there is an additional back chain field that points to the card previously punched so that error recovery for punch equipment checks can back up one card. The only exception to the format of Read/Write-TIC-Data is in buffers of files directed to the printer. In this case, immediate operation code CCWs (skips and spaces) are followed immediately by the next CCW.

#### Spool File Format

In addition to the data and CCWs contained in each SPOOL buffer, the first two doublewords contain forward and backward links to the next and previous buffers in the file. This two-way linkage allows the file to be backspaced/restarted from any point at any time. Also, it means that if I/O errors are encountered while reading one buffer, the file is put in system hold status. If purged, all buffers except those in error are released. The two-way chain allows this control of the file while preventing fragmentation by allowing pages to be assigned and released individually regardless of their ownership.

Each SPOOL file in the system is controlled by a Spool File Control Block (SFBLK) that is resident in storage. While the file is open, these blocks are chained from the devices (either real or virtual) that are processing the file, and from device type file anchors after the file is closed. There is one file chain each for printer, reader, and punch files. Each SFBLK contains information about the file that describes its owner and originator (these can be different for transferred files), the file name and type, and the class and number of copies for output files. All of these attributes can be examined and most can be changed by the file's owner or the system operator. The SFBLK also contains information such as the starting and ending buffer addresses for the file, the record size, certain file status flags, etc.

## SPOOL BUFFER MANAGEMENT

Real/Virtual Storage Management

Buffers used for the temporary storage of spool data on its way between DASD secondary storage and the user's virtual machine are allocated from a pool of virtual storage space that belongs to CP. This pool consists of the second 256K of virtual storage associated with the VMBLOK that controls CP's paging activities. This pool can be enlarged as a system generation option. Virtual storage buffers are allocated in one page increments by DMKPGT at the time the spool file is opened for either input or output. If no virtual storage space is available, the user is placed in a wait state until a buffer is freed by another user closing a file. This places limit on the number of concurrent spooling operations permitted by the system because spooling operates as a high priority task.

Real main storage is not allocated for a spooling buffer until a virtual machine actually issues a SIO that attempts to transfer data between the buffer and the user's virtual storage space. At this time, a page of real main storage is allocated to the buffer via the main storage paging manager. The buffer is locked in main storage (that is, is unavailable to be paged out) only for the amount of time necessary to transfer the data. After the data transfer is complete, the buffer is treated as a normal page of virtual storage, and can be selected to be paged out. This ensures that low usage spool files do not have buffers in real main storage, while the buffers for high usage files should remain resident. The Virtual Spooling Executive is insensitive to the location of the spool buffer in real storage, since all references to the data therein are accomplished through the dynamic address translation feature of the CPU.

DASD Space Allocation

While a spool buffer is active, it resides in real main storage or on the paging device. After it has been filled with data from the virtual machine or a real input reader, it is written to a page of secondary DASD storage. The allocation of pages on the

spooling disk(s) is managed by DMKPGT which is used to handle requests for both pages of virtual storage and semipermanent spool file residence. DMKPGT maintains separate allocation block chains for virtual storage and spooling pages. Each block contains control information and a bit map used to allocate pages on a single cylinder. If none of the cylinders allocated have any available pages, DMKPGT enters its cylinder allocation routine.

DMKPGT attempts to even out the spooling/paging I/O load by allocating cylinders in round robin fashion across channels and devices. In order to minimize seek times on a given device, an attempt is made to allocate cylinders as close to the relative center of the spooling/paging area as is possible.

Paging Device Support: All actual input/output for the page buffers on any device is controlled by the Paging I/O Executive DMKPAGIO and is discussed in that section in this publication.

## VIRTUAL SPOOLING MANAGER (DMKVSP)

The two functions of the virtual spooling manager are to simulate the operation of all spooled unit-record devices attached to the user's virtual machine, and to read and write the spool files associated with those devices. The following virtual devices are supported for spooling, with the exceptions noted:

- The IBM 2540 Card Reader/Punch, except for punch feed read and column binary
- The IBM 1403 Printer Models 2 and N1 (132 positions)
- The IBM 3211 Printer (150 print positions)
- The IBM 3505 Card Reader (except for mark senses reading)
- The IBM 3525 Punch (except for the card read, print, and data protect features).

The following consoles and terminals are also supported for spooling when entered into the directory as the

virtual system console:

- IBM 1052 Printer-Keyboard, Model 7 (via the 2150 Console)
- IBM 3210 Console Printer-Keyboard, Models 1 and 2
- IBM 3215 Console Printer-Keyboard, Model 1

It is assumed that all virtual printers have the universal character set feature. No checking is done on the spooled printer data. However, any UCS buffer commands issued by the virtual machine (load UCS buffer, block data checks, etc.) are ignored. It is up to the user and the installation to ensure that the output is directed to the proper real printer via use of the output CLASS feature described below. For the 3211 printer, Forms Control Buffer (FCB) commands are accepted and simulated by means of a virtual FCB maintained by the executive. The use of the virtual FCB is the only way to simulate end-of-form conditions reflected by the detection of a channel 9 or 12 punch. When the spooled file is directed to a real 3211 or 1403, the operator is responsible for loading the FCB or mounting the proper carriage tape.

If any of the unsupported unit-record features are required, they may be used by attaching the real device directly to the user's virtual machine. Thus, a 3505 reader could be used for the most part as a spooling input reader, but attached directly to a batch virtual machine when it is necessary to read mark sense cards.

#### Output File Processing

DMKVSP receives control from the Virtual I/O Executive DMKVIO when the user issues a SIO to a spooled unit record device. DMKVIO does not pass control until it has been determined that the device is available (that is, non-busy and with no interrupts pending). DMKVSP first determines if the device is currently processing a file. If it is, processing continues. If this is the first command issued by the given device, a new output file must be opened. An open subroutine is called to build the control blocks necessary to manage the file and to obtain virtual storage and DASD buffer space. Control is then returned to DMKVSP.

DMSVSP then analyzes and interprets the channel program

associated with the user's SIO. Each CCW is tested for validity of command, address, flags, alignment, protection, etc., and if the CCW is valid, the user's data is moved from his own virtual storage space to the buffer in the spooling virtual storage. When this buffer is full, it is written to a page of DASD secondary storage and a new buffer is obtained. The interpretation of the user's channel program continues until there are no more CCWs or until an error condition is detected which prohibits further processing. In either case, the device is marked as having the proper interrupts pending, a CSW is constructed, and DMKVSP exits to the main dispatcher. In contrast to nonspooled I/O, the user has remained in a pseudo-wait (IOWAIT) for the time it took to interpret the entire channel program.

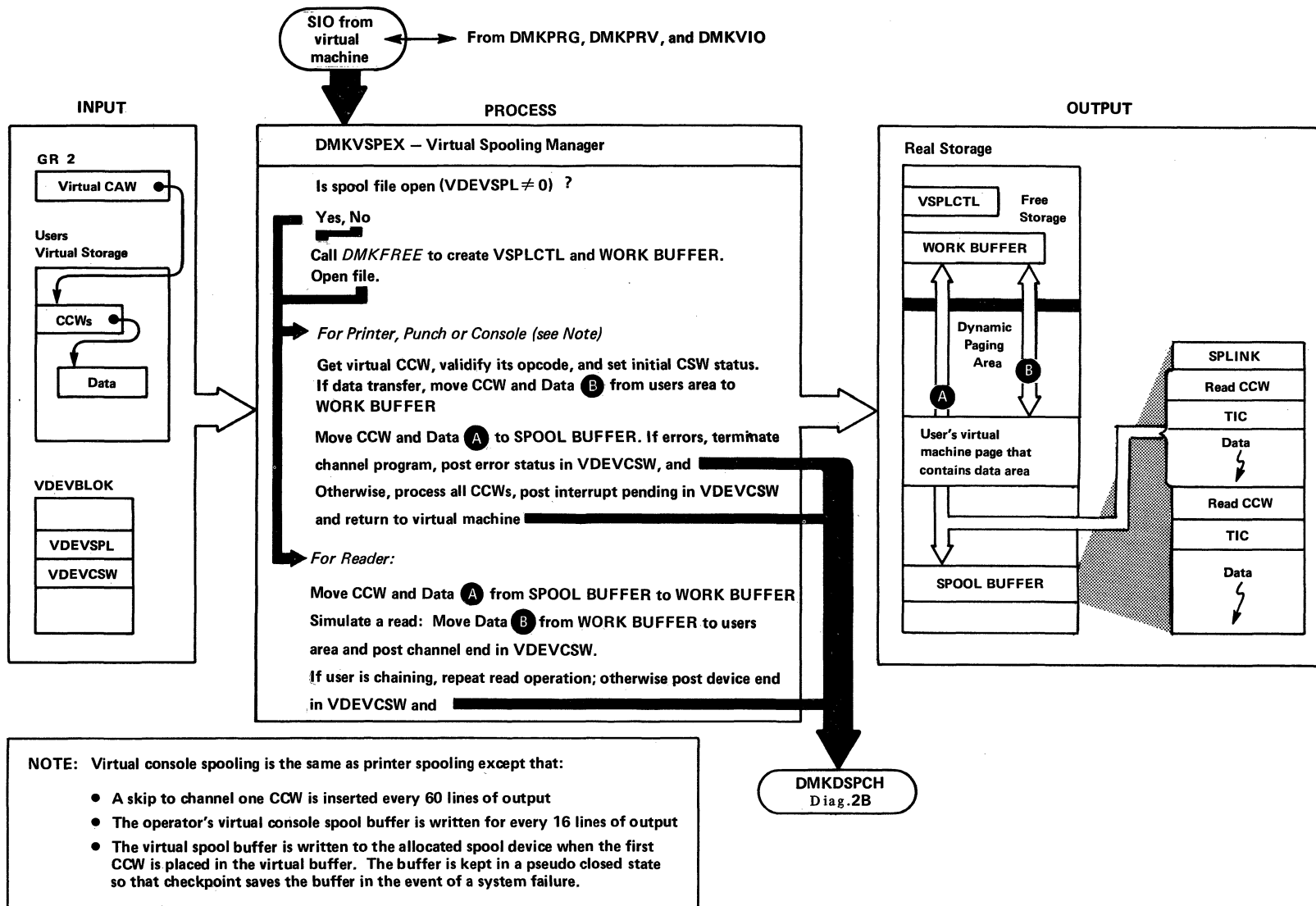
The output file can be logically closed by the user either by issuing an invalid CCW command code, or via the CP console function CLOSE. In either case, the device is cleared of pending interrupts, the file chains are completed, and the file is either queued for output on a real device of the proper type (printer or punch), or, if XPER is in effect, is queued for input to another user.

#### Input File Processing

Input file processing is similar to output file processing, except for the open and close functions, and the analysis of CCW commands and the direction of data movement. Many common routines are utilized to locate and verify CCWs, obtain buffer space, and to move the spooling data.

The difference in the open function is that instead of creating a new file, it is necessary to locate a reader file that already exists in the system. To do this, the open subroutine scans the SFBLKs chained from the anchor READERS in order to find a file with an owner userid that matches that of the caller. If a file is not found, a unit check/intervention required condition is reflected to the user; otherwise, its SFBLK is chained to the control block for the reader and the channel program is interpreted in the same manner as for an output file.

Diag. 4B1. Virtual Spooling Manager



After the input file is exhausted, a unit exception is reflected to the user machine, unless the user has requested either continuous spooling or that an EOF not be reflected. With continuous spooling, the unit exception is not reflected until the last file for that user is processed. If NOEOF is specified, the simulation terminates with a unit check/intervention required condition (similar to what happens if the EOF button on a real reader is not pushed).

In either case, the input file is then deleted from the system, unless the user has specifically requested that his input files be saved. If the file is saved, it can be re-read any number of times.

### Virtual Console Spooling

Support of the virtual console input and output is provided as an option of the VM/370 spooling capabilities. This support fulfills the following requirements:

- Provides hardcopy support for CMS Batch virtual machines.
- Allows DISCONNECTED virtual machines to spool virtual console output to disk instead of losing the output.
- Improves the performance of virtual machines that currently produce a large amount of console output.

Whenever a SIO IS ISSUED TO A VIRTUAL MACHINE CONSOLE| THE Virtual Console Manager (DMKVCN) determines if the spooling option is active. If it is, control is passed to the Virtual Spooling Manager at DMKVSPBP to insert the data into a spool file buffer. While console spooling utilizes, basically, the same code as printer spooling, the following exceptions are made:

- A skip to channel one CCW is inserted every 60 lines of output.
- The operator's virtual console spool buffer is written out every 16 lines of output.

- The virtual space buffer is written out to the allocated spool device when the first CCW is placed in that virtual buffer. The linkage area of the virtual spool buffer takes the form of a CLOSE file to allow checkpoint (DMKCKP) to recover the active spool file in the event of a shutdown due to system failure. The data in the virtual buffer, not yet written out to the spool device will not be recovered.

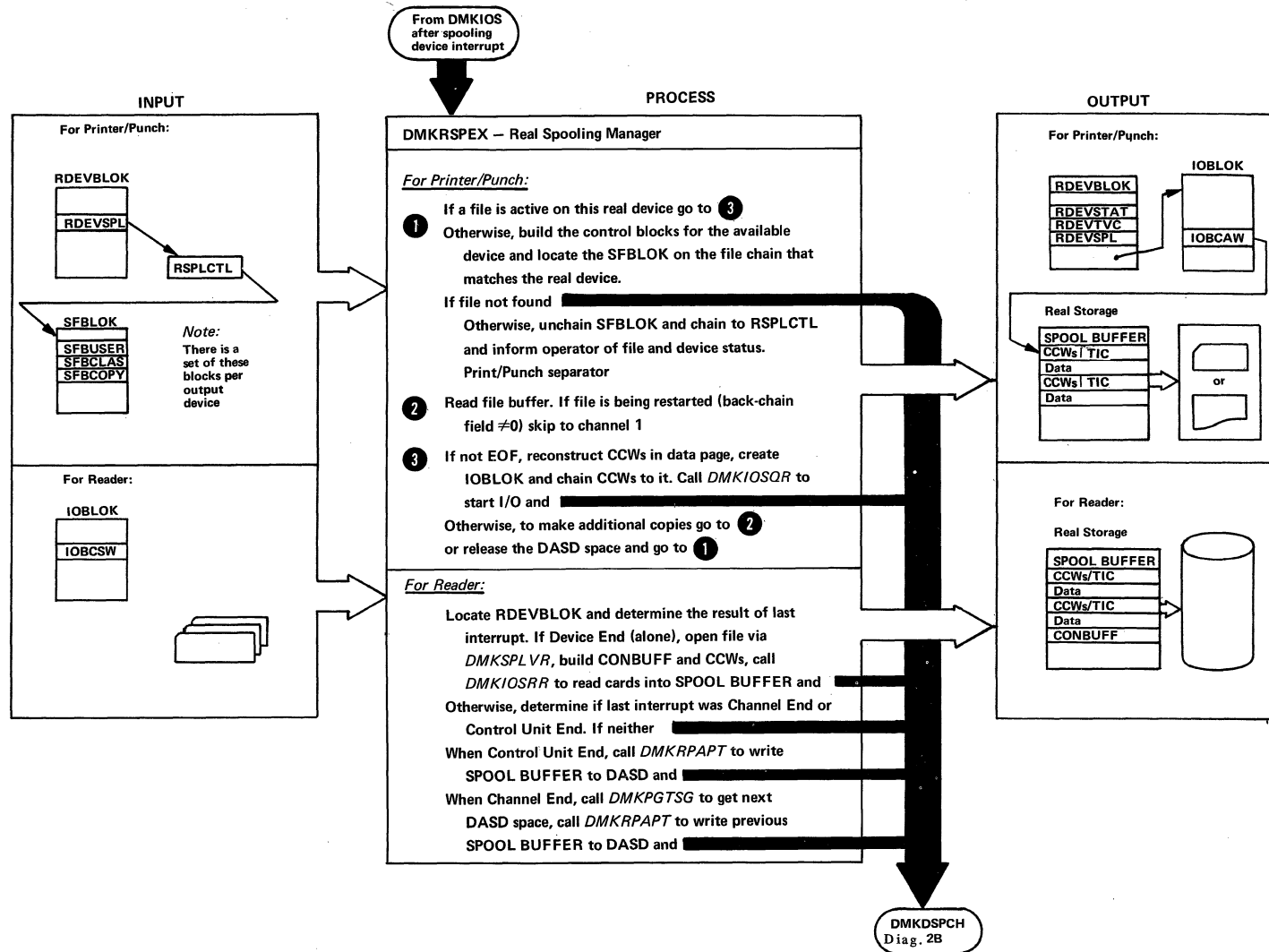
To maintain a pseudo closed file status for console spool files, DMKSPL now assigns spool identifications to all output spool files where they are first queued.

A virtual system reset, device reset, or IPL will not close the virtual console spool file. The LOGOFF, FORCE, or DETACH of virtual console commands will close the virtual console spool file. The SHUTDOWN command will close the operator's console spool file. If the SHUTDOWN command is issued by a Class A user other than the operator, the console spool file for both the user and operator will be closed.

### REAL SPOOLING MANAGER (DMKRSP)

The real spooling manager operates the real unit record devices that are attached to the system and that are used to transcribe input data into reader spool files and user output spool files onto the real printers and punches. The executive optimizes the use of main storage and the CPU rather than running the system unit record devices at their rated speeds. DASD input files are not double buffered and under periods of peak load, input and output devices tend to run in bursts. However, command chaining is used for all unit record channel programs so that the devices are running at their maximum speed with a minimum of interruptions.

Diag. 4B2. Real Spooling Manager





## Output File Processing

Both the input and output functions of DMKRSP are interrupt driven. Thus, DMKRSP does not process unless an internally or externally generated not-ready to ready device end interrupt occurs. External interrupts are generated by the hardware in the normal manner, while internal, "psuedo interrupts," are generated by the software when an output file has been queued on the real printers or punches file chain, or when the operator issues a START command to a drained device.

Upon receipt of the initial device end for a printer or punch, DMKRSP searches the appropriate file chain for the SFBLOK of a file whose class matches that of the device that was made ready. When the SFBLOK is located (provided the file is not in a HOLD status), it is unchained from the output queue and chained to the real device block that services the file. A page of real main storage is then obtained for use as a buffer, and the output separator routine (DMKSEP) is called to print output identifier pages. When DMKSEP returns control to DMKRSP, the first buffer of the file is paged into real main storage, and the CCWs in the channel program that it contains are adjusted so that their data addresses correspond to the real addresses at which the data resides. The real SIO supervisor (DMKIOSQR) is then called to start the channel program, and DMKRSP exits to the dispatcher (DMKDSPCH) to await the interrupt.

When the channel end/device end interrupt for the completed buffer is unstacked to DMKRSP, the forward chain file link field is used to locate the next buffer. This buffer is paged-in, and the process is repeated until the final buffer is processed. At this point, the number of copies requested for the file is decremented. If the number of copies is 0, processing is terminated and the file is deleted from the system; otherwise, the process is repeated as many times as is necessary.

When file processing is complete, a scan of the appropriate output queue is again made, and if a file is found it is processed. If the queue is empty, or if a file with a matching class is not found, an exit is taken to DMKDSPCH to wait for another ready interrupt.

Output file processing can be modified by either the system operator via the spooling support command or as a result of system errors. The operator commands allow a given file to be backspaced or restarted, and the files of individual users or the whole system to be held and released for output in a very flexible manner. I/O errors also affect the spooling system, and a description of how they are processed is in the section "Error Recovery."

## Input File Processing

Reader file processing is initiated by the receipt of a device end interrupt from a spooling card reader. No explicit operator command is required to start the processing of an input file. When the device end is unstacked to DMKRSP, an open subroutine is called to build the necessary control blocks and to obtain the virtual, real, and DASD buffer space required for the file. A channel program to read 41 cards is built in the buffer, and DMKIOSQR is called to start the reader.

When the interrupt for the first buffer is unstacked, the first card is checked for its validity as a userid card. The minimum information that this card must contain is the userid of the owner of the input file. It may appear anywhere on the card, with the restriction that it must be the first information punched. Optional information on the userid card can include a file name and type and/or the class of the virtual card reader to which the file is to be directed. If the userid is valid, the file processing continues; otherwise, the operator receives an error message and processing is terminated.

After each file buffer is read, it is written onto disk by the paging I/O routines in the same manner that virtual output files are handled. When a unit exception signaling physical end of file is received from the reader, the file is closed by writing the final buffer to disk and completing and queueing the SFBLOK to the readers file chain. If the owner of the file is currently logged in, he is given a message indicating that a file has been read and the appropriate card reader is posted with a device end interrupt.

Accounting Card Processing

Various routines in CP accumulate, format, and punch account cards that contain system usage information for certain users. These routines format the information into an 80-column card image preceded by a punch CCW and call DMKACOAQ to queue the card for real output. DMKACOAQ calls DMKACOPU to punch the card on a real punch if one is available; otherwise, the card is queued in main storage until a punch is free. When a punch finishes processing its last file, a test is made to see if any accounting cards have been queued. If they have, DMKACOPU is called to process them.

In addition to the cards generated by CP to account for a virtual machine's use of system resources, the user may request cards to be punched in order to account for the use of virtual machine resources by jobs running under his userid. In order to do so, the user must have specified the account option (ACCT) when initially entered into the directory.

In order to punch an accounting card, the user must issue a X'004C' diagnose instruction with a pointer to a parameter list containing the "charge to" information. If the pointer is zero, the accounting card will be punched and will contain the user's own identification taken from his VMBLOK.

When the user accounting option is being utilized, the user must keep in mind that each additional accounting record requested is occupying real storage space. Degradation of system performance will occur if available storage becomes filled with accounting data.

SPOOLING COMMAND SYSTEM

The spooling command system provides an interface between the user, the system operator, and the spooling system itself. There are three types of spooling commands.

- Those that affect virtual devices

- Those that affect real devices
- Those that affect SPOOL files that are queued within the system

The commands that affect virtual devices are generally available to all system users, and a user can only affect the status of devices that are attached to his own virtual machine. Commands that affect the status of the real system's spooling devices are restricted to use by the system operator. Commands that affect closed spool files that are awaiting processing are generally available to all users, with some additional capabilities assigned to the system operator. For example, a user may alter the characteristics only of those files that have an owner's userid that matches his own, whereas the system operator may change any SPOOL file in the system.

File States and Attributes

Each spool file in the system has a number of attributes that are assigned to it, either explicitly or by default, at the time that it is created. These attributes and their values are as follows:

- Filename and type can be 24 character fields. Either or both can be replaced by a user-supplied value.
- Spoolid number is a system-assigned number between 1 and 9999. It is automatically assigned when the file is created (input) or closed (output), and is unique within the system. The file's owner, the device type, and the id number are specified. Usually, the userid defaults to the identification of the user issuing the given command. Since the identification number rather than the file name and type is used as an identifier, duplicate user-assigned names do not present an identification problem.
- The number of logical records (cards or print lines) in the file is an integer between 1 and 16 million. For printer files, the record count also includes any immediate operation code space or skip CCWs.

- The Originating User is the identification of the files creator, if the file has been internally transferred from the originator's printer or punch to the new owner's card reader.
- The number of copies requested for an output file is a number between 1 and 99. Unless altered by the user or operator, it defaults to 1.
- The device type is used by DIAGNOSE for a file transferred to a reader to determine the virtual type of output device.

In addition to those attributes, a file that is queued for real output or virtual input always has a class associated with it. A class is a single alphameric character from A through Z or from 0 to 9. It is used to control both the real or virtual device on which the file will be printed, punched, or read, and the relative priority and sequence of output on the device. While each file is assigned a single class, each real spooling output device be assigned from one to four classes. The device then processes only files that have a class attribute that corresponds to one of its own, and will process these files in the order that its own classes are specified.

For example, if a printer is assigned the classes A, D, 2, it processes any printer file with a class of A before it searches the printer output queue for a file with class D. All class D files are printed before class 2 files.

The output class for a file is assigned at the time the file is created and is the class that is associated with the virtual device that created it. While each real spooling device can have up to four classes, each virtual spooling device can have only one. When a user logs onto to the system, the class associated with a device is the one defined in his directory entry for that device. However, he can alter this class at any time via the spool command. As files are created and closed by a device, they take on the device's output class.

After they are closed and are awaiting output, their class can be changed via a CHANGE command issued either by the file's owner or the system operator. The system operator can alter the system generated output class(es) of a real output device via the START command.

Output files transferred to a user's virtual reader can also be controlled by class. If the receiving user has several readers, the input to each can be limited to files of a certain class. In addition, the ORDER command allows sequencing of input files by class as well as spoolid number.

Output priorities can also be managed by altering the hold status of a file. Individual users can alter the hold status with the CHANGE command, while the system operator can change (hold or free) the files of specific individual users.

#### Virtual Device Spooling Commands

These commands affect the status of a user's virtual spooling devices:

<u>Command</u>	<u>Meaning</u>
CLOSE	Terminates spooling operations on a specified device. It clears the device of any pending interrupt conditions, and for output files completes and queues the file for real output. Optional parameters allow the user to specify a filename and type, and to override for the given file any standard class, hold/nohold or copy parameters set into the output device by the spool command.
SPOOL	Establishes the file attributes that apply to files created on, or read by, the given device. It establishes the CLASS that will be in effect, whether: files are to be automatically held, input files are to be saved or purged after reading, and output files are to be directed to the real system printers and punches or are to be transferred to a user's virtual reader.



### Real Device Spooling Commands

The operator can use these commands to control the activity of the real spooling devices:

<u>Command</u>	<u>Meaning</u>
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BACKSPAC	Backspaces an active spooling device for either a specified number of pages (printers only) or to the beginning of the file (printers or punches).
DRAIN	Stops the operation of a specified output or input device after it has finished processing the file on which it is currently working. A printer must be drained prior to the issuance of the LCADEUF command. Unit record devices are normally drained prior to system shutdown.
START	Restart a device after it has been drained. Optional parameters allow the operator to specify the spooling output class for the output device, and if output separator records will be created.
FLUSH	Immediately halt the output on the specified device and either flush that copy of the file from the system, or put it into the system hold status for future processing.
REPEAT	Supplement the number of copies requested by the user for the file when it was created. The operator can specify a number from 1 to 99 that is added to the number specified by the user.
LOADEUF	Load the Universal Character Set Buffer of the FCB of the specified printer with the specified image. If requested, the system verifies the loading by printing its contents on the affected printer.
SPACE	Force the output on the specified printer to be single spaced, regardless of the skipping or spacing commands specified by the file's creator.

Spool File Management Commands: The spooling commands are used to alter the attributes and status of closed spool files that are queued and awaiting processing. When a command applies to an individual file, the device type (RDR, PUN, PRT) and the spoolid number must be provided in order to identify the file. It should be noted that in most commands requiring a spoolid, the keyword class followed by a valid spool class or the keyword ALL are acceptable substitutes for the spoolid number. This causes the command to be executed for all files of the given class or device type. The userid is assumed to be the identification of the user issuing the command, except that the system operator must explicitly supply the identification of the user whose files he wishes to affect or he must specify the keyword SYSTEM which gives access to all files (valid for CHANGE, PURGE, ORDER, and TRANSFER commands also).

<u>Command</u>	<u>Meaning</u>
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CHANGE	Change the filename and type, the number of copies, and the class of the specified file. Any of the above attributes of a file can be determined via the QUERY command.
HOLD	Place, via the system operator, the specified file in a hold status. The file will not be printed or punched until it is released by the system operator. The operator can hold any user files by device type.
FREE	Opposite of the HOLD command. Allows a file or group of files that were previously held to become available for processing. However, the user cannot reset a hold set by the operator via the HOLD command.
PURGE	Removes unwanted spool files from the system before they are printed or punched.
ORDER	Reorder the input files in a virtual card reader. It can order files by identification number, by class, or by any mixture of the two.
TRANSFER	Transfer a virtual input to another user's virtual reader without any processing. The TRANSFER command causes a changing in the owning userid field in the file's SFBLOK.

## SPOOLING ERROR RECOVERY

Unit Record I/O Errors

I/O errors on real spooling unit record devices are handled by a transient routine that is called by DMKIOS after it has sensed the unit check associated with the error on a spooling device. If appropriate, a restart CAW is calculated and DMKIOS is requested to retry the operation, in some cases waiting for a device end that signals that the failing device has been made ready after manual corrective measures have been taken. If after retrying the operation the error is unrecoverable, DMKIOS is informed that a fatal error has occurred. DMKIOS then unstacks the interrupt, flagged as a fatal error, and passes control to real spooling executive. The routines that handle unstacked interrupts in real spooling executive only see operations that have been completed correctly or those that are fatal errors. If a fatal error is unstacked, the recovery mechanism depends on the operation in progress.

For fatal reader errors, processing of the current file is terminated and any portion of the file that has been read and stored on disk is purged. The file's owner is not informed of the presence of a fractional part of the file in the system.

For fatal printer or punch errors, the SFBLOK for the partially completed file is re-queued to the appropriate output list and processing can be resumed by another available printer or punch, or can be deferred until the failing device is repaired.

In any case, the failing device is marked logically offline, and no attempt is made by the system to use it until the operator varies it back online via the VARY command.

DASD Errors During Spooling

DASD I/O errors for page writes are transparent to the user. A new page for the buffer is assigned, the file linkage pointers are adjusted, and the buffer is

rewritten. The failing page is not de-allocated and no subsequent request for page space granted access to the failing page. If an unrecoverable error is encountered while reading a page, processing depends on the routine that is reading the file. If the processing is being done for a virtual reader, the user is informed of the error and a unit check/intervention required condition is reflected to the reader. If the processing is being done for a real printer or punch, the failing buffer is put into the system hold status, and processing continues with the next file. In either case, the DASD page is not de-allocated and it is not available for the use of other tasks.

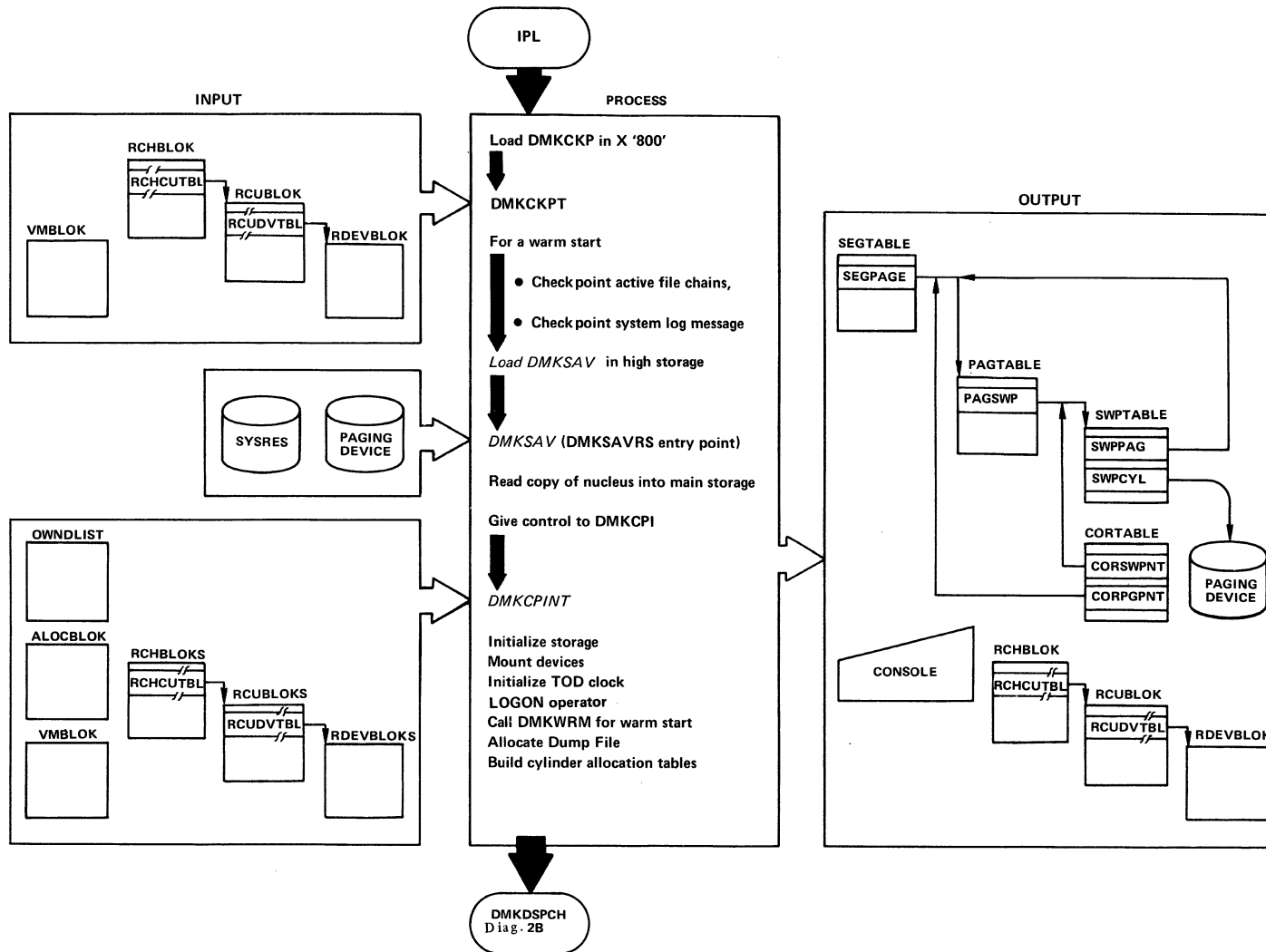
DASD Spool Space Exhausted

If the space allocated for paging and spooling on the system's DASD volumes is exhausted and more is requested by a virtual spooling function, the user receives a message and a unit check intervention required condition is reflected to the virtual output device that is requesting the space, the output file is automatically closed and it is available for future processing. The user can clear the unit check and retry the operation periodically in the hope that space is free or completely restart later from the beginning of the job. If the task requesting the space is the real spooling reader task, the operator receives an error message and the partially complete file is purged. Any time the spooling space is exhausted, the operator is warned by a console message and alarm. However, the system attempts to continue normal operation.

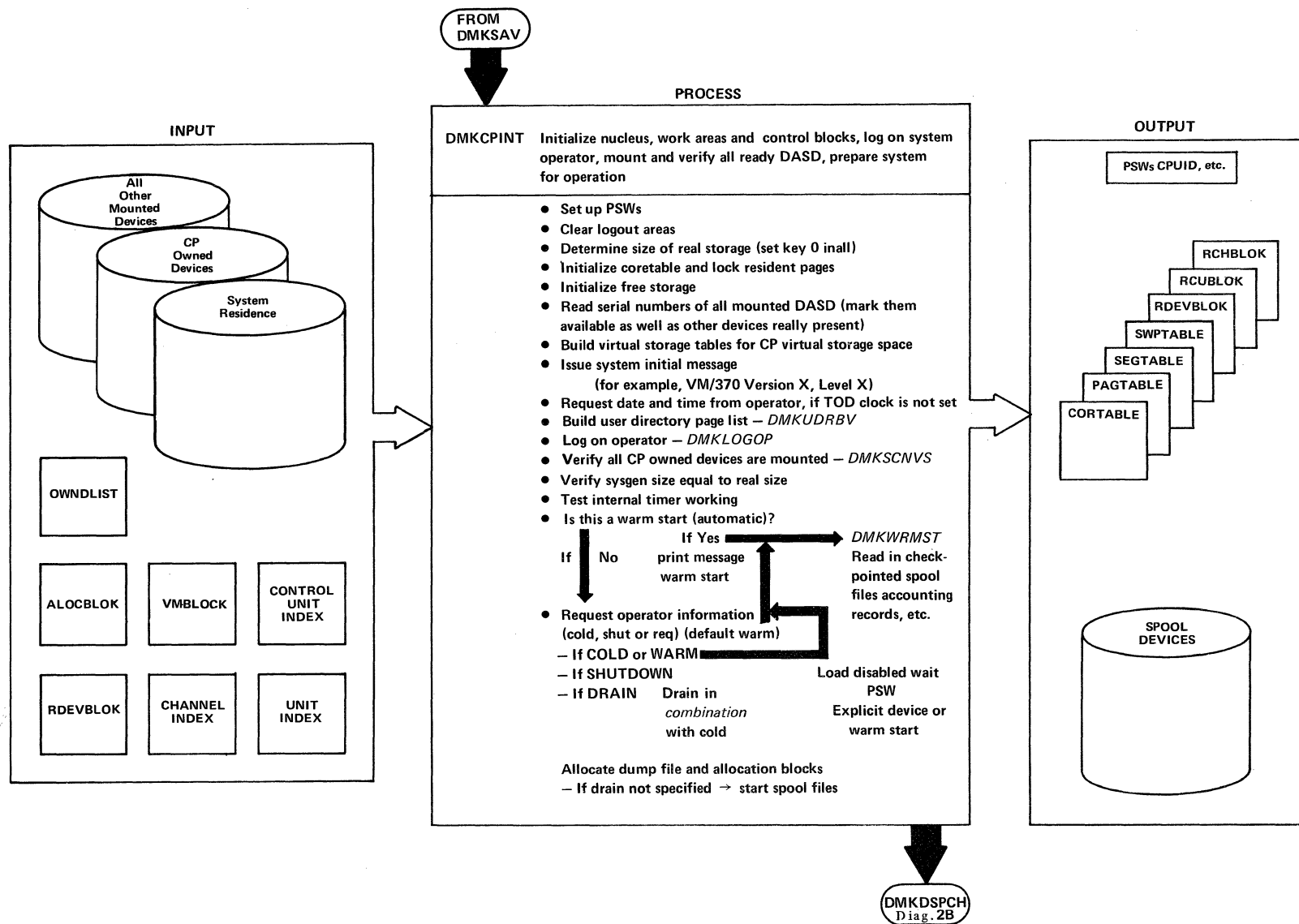
CP INITIALIZATION

System initialization starts when the operator selects the DASD device address of the VM/370 control program System's Residence Volume (SYSRES) and presses the IPL button. The System/370 hardware reads 24 bytes from record 1 of cylinder 0 on SYSRES into location 0 of main storage. This record consists of an initial PSW and a channel program. The channel program is used to read the module DMKCKP into location X'800' and give it

Diag. 5B. CP Initialization

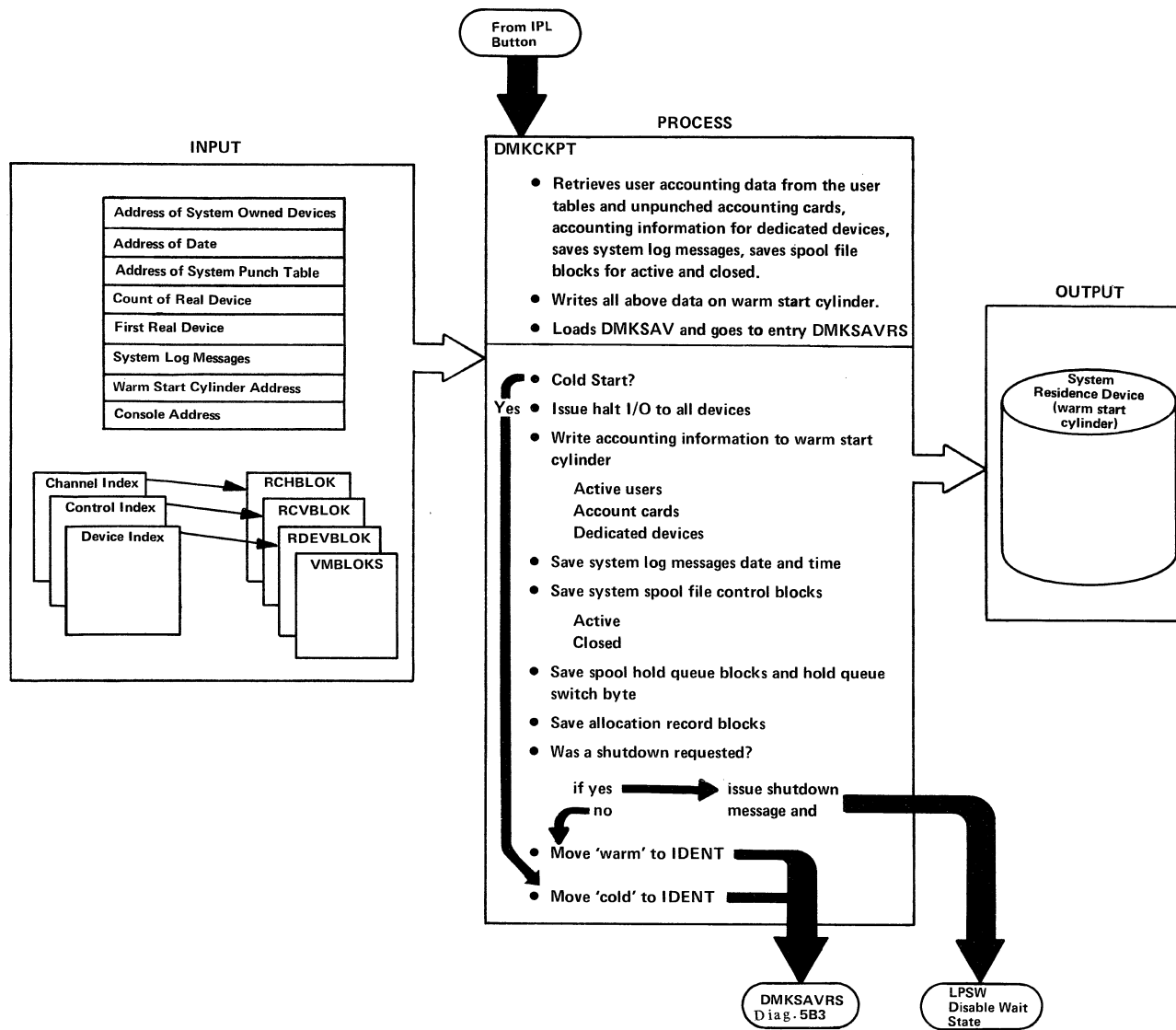


| Diag. 5B1. CP IPL

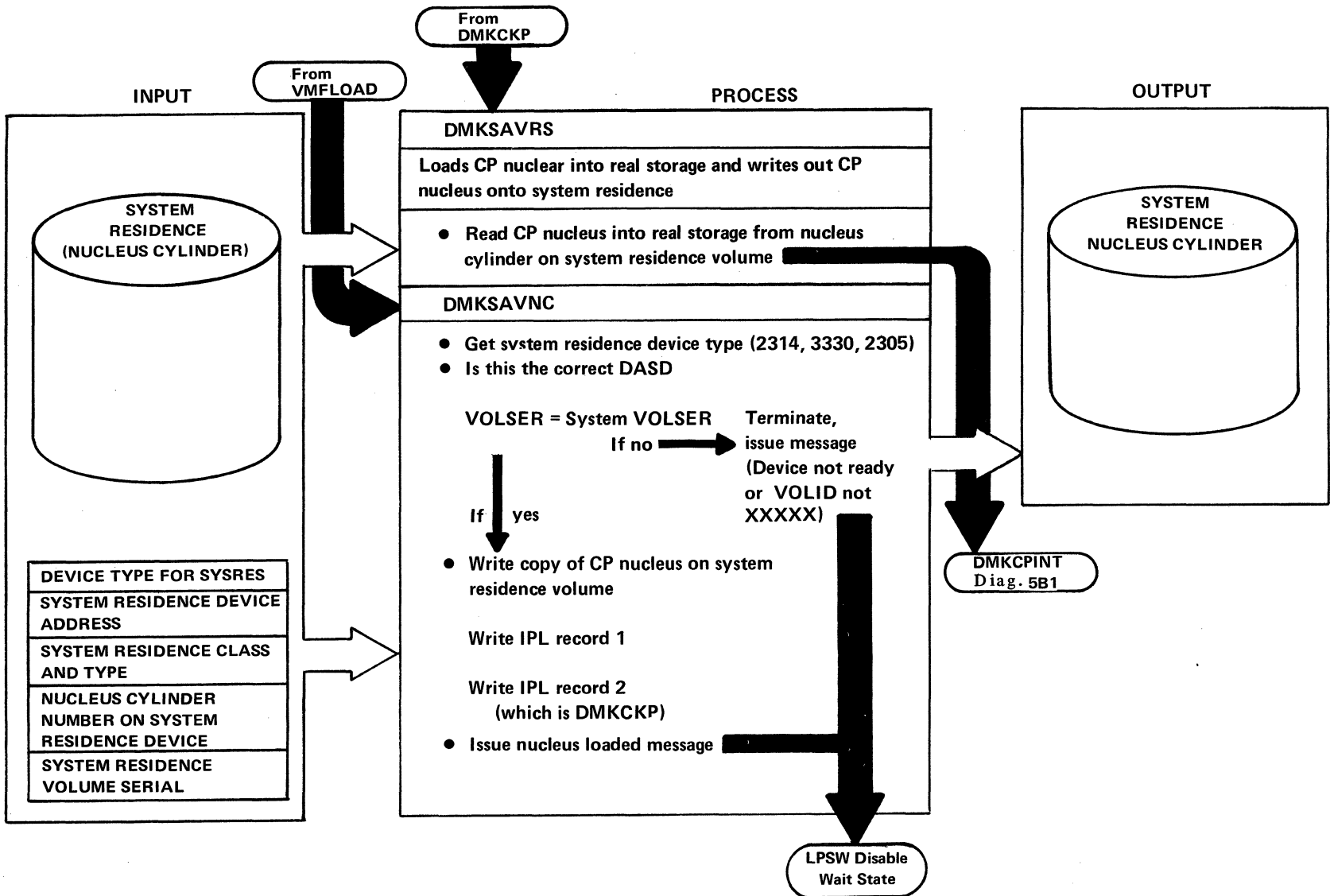




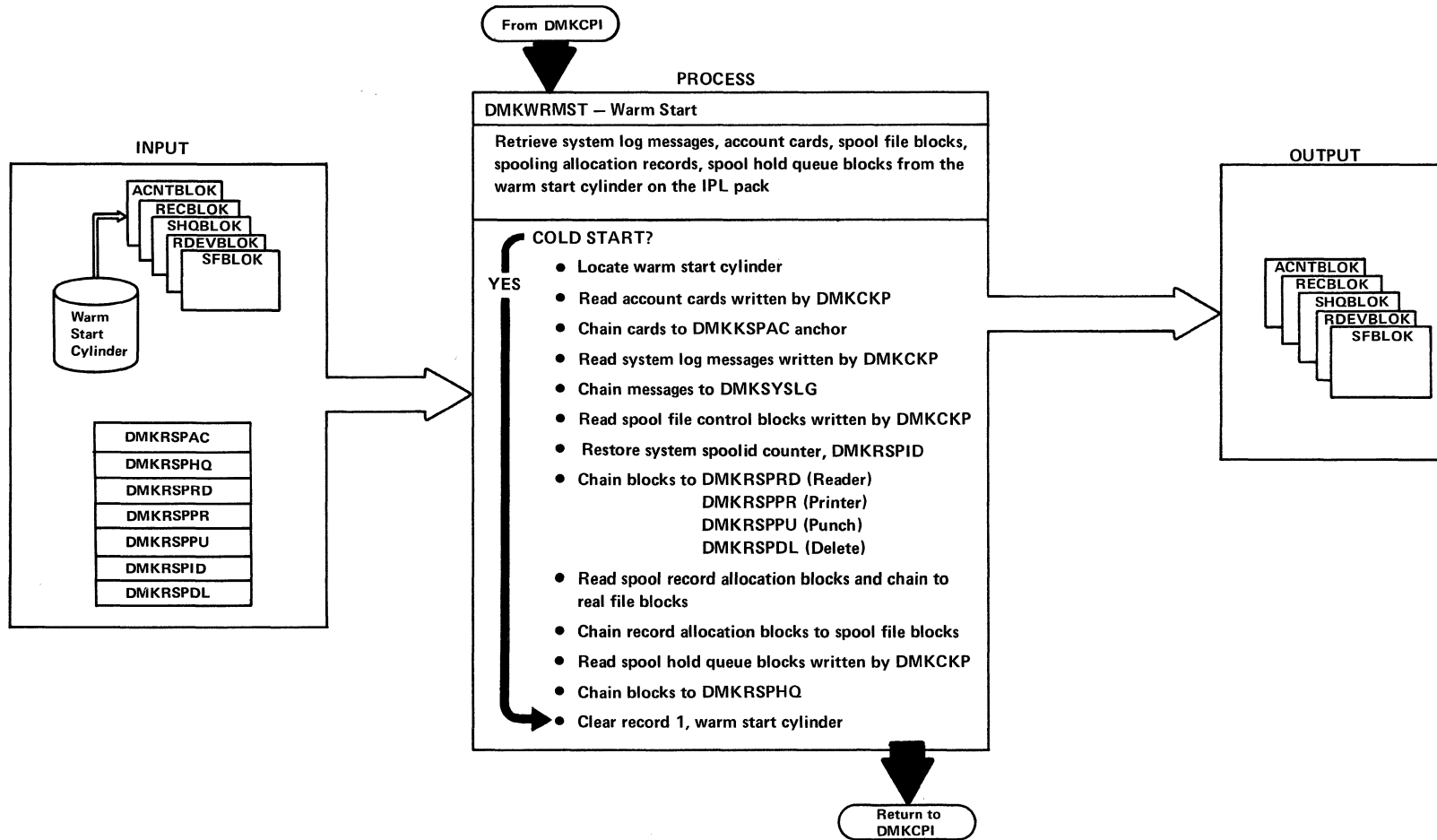
Diag. 5B2. Check Point



Diag. 5B3. Save System



Diag. 5B4. Warm Start



control. DMKCKP checks location CPID in module DMKPSA. If this location contains the value CPCP or WARM, then DMKCKP checkpoints the active file chains and saves the system log messages and accounting information; otherwise, a cold start is performed and the checkpointing is not done.

If location CPID contained the value CPCP, then checkpoint loads a wait state PSW at this time.

If location CPID does not contain the value CPCP, then DMKCP loads DMKSAV and passes control to it at entry point DMKSAVRS. DMKSAV reloads a page image copy of the CP nucleus into real storage starting at page 0. When DMKSAV is finished, control is transferred to DMKCPI. DMKCPI performs the main initialization function. This includes calling DMKWARM to perform the warm start function. When DMKCPI has finished it, it passes control to DMKDSPCH. DMKDSPCH loads a wait state PSW to wait for work.

#### FREE STORAGE MANAGEMENT

DMKFRE is responsible for the management of free storage, and is used within the control program for obtaining free storage for I/O tasks, CCW strings, various I/O buffers, etc. It is used, in fact, for practically all such applications except real channel, control-unit, and device-blocks, and the CORTABLE.

Block sizes of 30 doublewords or less, constituting about 99 per cent of all calls for free storage, are grouped into 10 subpool sizes (3 doublewords each), and are handled by LIFO (push-down stack) logic. Blocks of greater than 30 doublewords are strung off a chained list in the classic manner.

Subpool blocks are generally obtained, when none are available, from the first larger sized block at the end of available free storage with the smaller sizes. Large blocks, on the other hand, are obtained from the high-numbered end of the last larger block. This procedure tends to keep the volatile small subpool blocks separated from the large blocks, some of which stay in storage for much longer periods of time; thus, undue fragmenting of available storage is avoided.

DMKFRE initially starts without any subpool blocks; they are obtained from DMKFREE and returned to DMKFRET on a demand basis.

The various cases of calls to DMKFREE for obtaining free storage, or to DMKFRET for returning it, for subpool sizes and large sizes, are handled as follows:

#### Call to DMKFREE for a Subpool Size

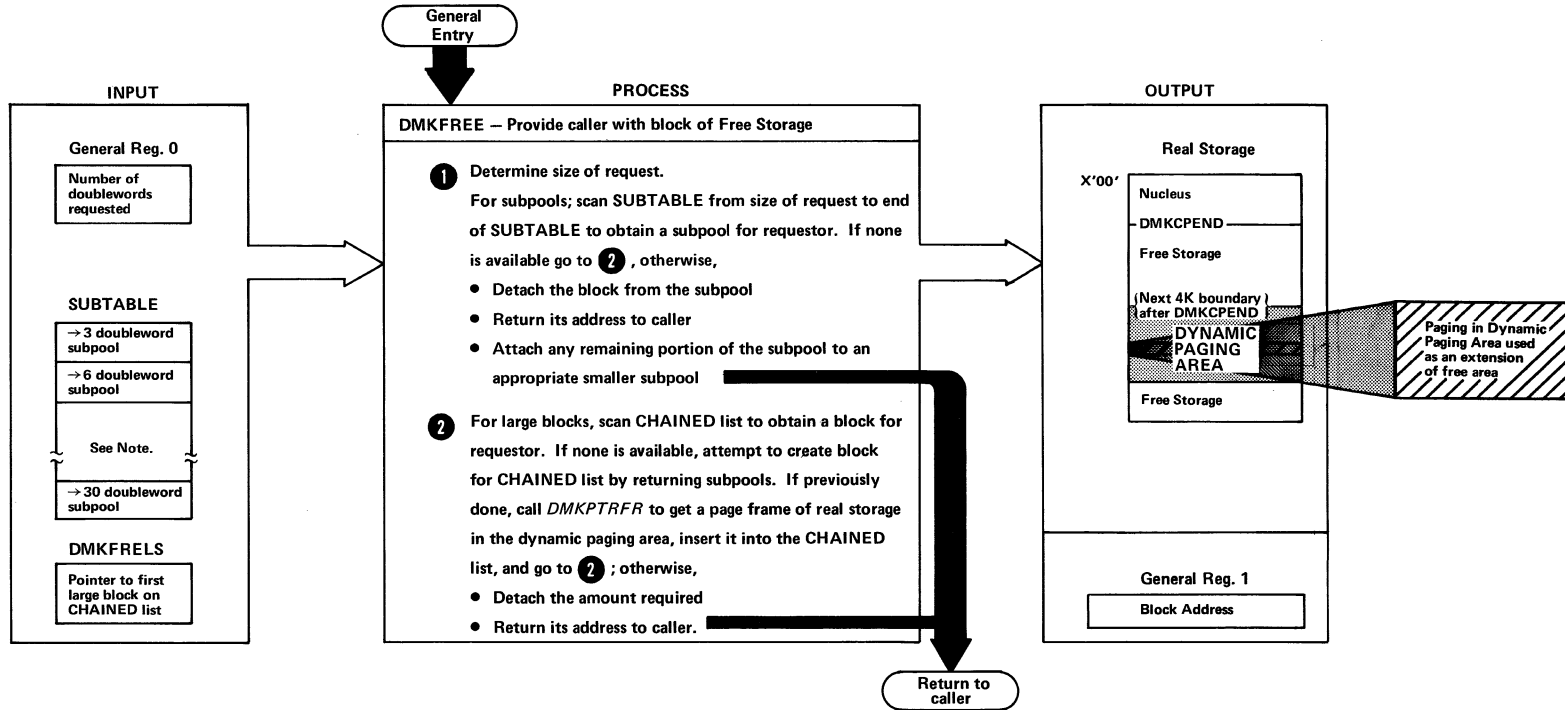
Subpool Available: If a call for a subpool size is made and a block of the suitable size is available, the block found is detached from the chain, the chain patched to the next subpool block of the same size (if any), and the given block returned to the caller.

Subpool Not Available: If there is no suitable block when a call to DMKFREE is made for a subpool size, a check is made to see if any larger subpool block can profitably be split up into the size requested and another subpool size. If this is feasible, the larger block is detached from its subpool and split. The requested block is returned to the caller, and the remaining block is attached to its subpool. Otherwise, the chained list of free storage is searched for a block of equal or larger size. The first block of larger or equal storage is used to satisfy the call (an equal-size block taking priority), except that blocks within the dynamic paging area are avoided if at all possible. If no equal or larger block is found, all the subpool blocks currently not in use are returned to the main free storage chain, and then the free storage chain is again searched for a block big enough to satisfy the call. If there still is no block big enough to satisfy the request, then DMKPTRFR is called to obtain another page of storage from the dynamic paging area, and the process is repeated to obtain the needed block.

#### Call to DMKFREE for a Large Block

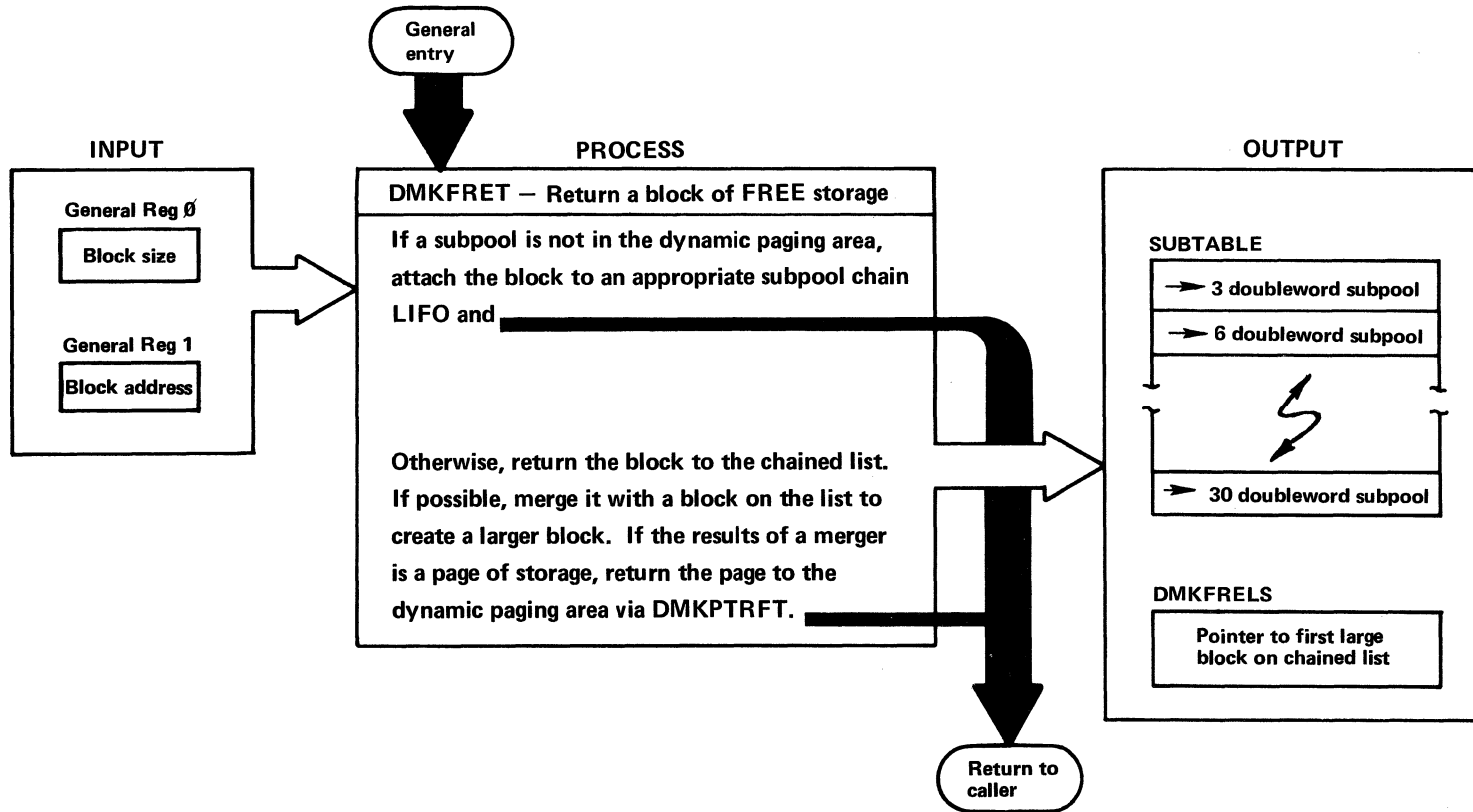
If a call to DMKFREE is made for a block larger than 30 doublewords, then the chained list of free storage is searched for a block of equal or larger size. If an

Diag. 6B1. Free (Provide) a Block of Storage



Note: Block sizes of 30 doublewords or less are grouped into 10 subpools (3, 6 . . . 30 doublewords, multiples of 3 doublewords); the subpools are processed LIFO. The SUBTABLE contains pointers to the subpools. Block sizes greater than 30 doublewords chained together by a CHAINED list. DMKFRELS points to the first block on the list.

Diag. 6B2. Return a Block of Storage



equal size block is found it is detached from the chain and given to the caller. If at least one larger block is found, the desired block size is split off the high numbered end of the last larger block found, and given to the caller. If no equal or larger block is found, DMKPTRFR is called to obtain another page of storage from the dynamic paging area, and the above process is repeated (as necessary) to obtain the needed block.

#### Call to DMKFRET for a Subpool Size

If a subpool size block is given back via a call to DMKFRET, the block is attached to the appropriate subpool chain on a LIFO (push-down stack) basis, and return is made to the caller. If, however, the block was in a page within the dynamic paging area, the block is returned to the regular free storage chain instead.

#### Call to DMKFRET for a Large Block

If a block larger than 30 doublewords is returned via DMKFRET, it is merged appropriately into the regular free storage chain. Then, unless the block was returned by DMKPRETR (see the section, "Initialization", a check is made to see if the area given back (after all merging has been done) is a page within the dynamic paging area. If so, it is returned to the dynamic paging area via DMKPTRFT for subsequent use.

#### Initialization

The number of pages allocated to free storage depends upon the number of storage boxes upon which the VM/370 control program is running, and is initialized by DMKCPINT (usually 6 pages per 256K). DMKPRETR is called by DMKCPINT to merge available blocks of storage into the regular free storage chain regardless of their size.

#### CONSOLE FUNCTIONS

DMKCFM analyzes VM/370 control program commands and pass control to the appropriate routine to handle the command. DMKCFM can be entered via the attention key at the user's terminal or directly from a virtual machine.

When a console interrupt occurs via the attention key at the user's terminal, DMKIOSIN calls DMKCNSIN to handle the unsolicited interrupt, then DMKCNSIN calls DMKCFMBK.

DMKCFMBK first calls DMKFREE to obtain storage for an 18 doubleword input buffer. Next, DMKQCNWT is called to send the message CP to the terminal to inform the user that he has entered console function mode. DMKQCNRD is then called to read the console function request.

DMKCFMEN is the entry point for commands coming directly from the virtual machine. DMKPRGIN enters here when a DIAGNOSE instruction with a code of 8 is detected. The address of an 18 doubleword input buffer is passed in register 1; therefore, a read to the terminal is not needed.

After either the read to the terminal or entry from the virtual machine, DMKSCNFD is called to find the command type. On return from DMKSCNFD, register 1 points to the start of the command and register 0 contains the length of the command. The entered command is matched against a list of valid commands. The list contains a 16-byte entry for each command. Each entry contains 8 bytes for the name, 2 bytes for class mask, 2 bytes for an abbreviation count, and 4 bytes containing the routine address. If the entered command matches an entry in the list, it is then checked to ensure that a valid abbreviation for the command has been used. If this test is not successful, DMKSCN continues to scan the list for a valid command. Should the abbreviation be valid, a check is then made to determine if this user is of the proper class to use the command entered. If this is successful, DMKCFM then calls the appropriate routine to process the command.

After the command has been processed, control is returned to DMKCFM. There are three possible returns. On a normal return, the input buffer is scanned to see if there are any more commands. If none exist, DMKCFM returns to the virtual machine (if entered via

DIAGNOSE) or calls DMKQCNRD to read the next command from the terminal. On a return plus 4, the VMCFWAIT bit is turned off to allow the virtual machine to run. DMKFRET is called to return the input buffer storage. Then control returns to either the virtual machine, if entered via a diagnose or to DMKDSPCH, if entered via the attention key. On a return plus 8, the operation is the same as plus 4 except the VMCFWAIT bit is left on.

## SYSTEM/USER INTERFACE

### Attaching a User to the System

After CP has been initialized, the communication lines are enabled by DMKCPVEN. Then an individual user is attached to the system using the following steps:

#### 1. Terminal Identification

When the Control Program receives the initial interrupt from a terminal on an enabled line (normally initiated by a user dialing in on a data-set), the DMKCNSID routine is entered. DMKCNSID determines the terminal device type, stores this information in the terminal device block, writes the online message and puts the terminal line in a state to receive an attention.

#### 2. Attention from User

After the online message has been typed at the user's terminal, and he has pressed the Attention key, DMKCNSIN (the console-interrupt routine) calls DMKBLDVM to build a skeleton VMBLOK for the user. At this time, the USERID is LOGONxxx, where xxx is the terminal real device address, and a flag is set to indicate that the user has not yet completed the LOGON process.

Then DMKCNSIN calls DMKCFMBK, which types a single blank at the terminal, issues a read to the terminal, for the user to enter his first command (normally LOGON or DIAL).

#### 3. First Command from User

After the first command has been entered by the user, DMKCNSIN further determines the type of terminal. If the terminal is a 2741, DMKTRMID is called to identify it as either a 2741P (PTTC/EBCD) or a 2741C (Correspondence) terminal. If successful, the correct device type and translate tables for input and output are set; if not, flags are set to indicate the terminal is not yet identified.

Then control is returned to DMKCFMBK, which determines if the first command is valid (for example, LOGON, MSG, or DIAL). If the first command is not valid, a restart message is given, and the read to the terminal posted again for the first command. If the first command was LOGON (or its abbreviation), DMKLOGON is called to complete the process of attaching the user to the system.

#### 4. LOGON of User

Operations performed by DMKLOGON include the following:

- Ensuring that the maximum number of users allowed on the system is not being exceeded.
- Obtaining the userid from the command line, and checking for a possible password and other optional parameters.
- Checking the userid and password (entered separately if not on the LOGON command line) against entries in CP's directory of users.
- Ensuring that the user is not logged on at another terminal (an error condition), or reconnecting the user if he was running, but in the disconnect mode.
- Obtaining pertinent information on the user's virtual machine from the User Machine Block portion of the directory.
- Storing the correct userid (replacing the LOGONxxx userid used up until now), virtual storage size, and other vital information in the user's VMBLOK.
- Allocating and initializing segment, page, and swap tables (necessary for handling of the user's virtual storage).



- Allocating an extended VMBLOK (ECBLOK) if the user's virtual machine has the capability of running in the extended control mode.
- Allocating and initializing virtual device blocks, control unit blocks, and channel blocks, using information from the User Device Blocks portion of the directory.
- Establishing links (as feasible) to all DASD devices included in the user's directory, the accessibility of any disk being determined by the user access mode in the user's directory, and whether any other user(s) are presently linked to the disk, in read-mode and/or write-mode.
- Initializing all other virtual device blocks as appropriate, such as reader, punch, printer, and terminal.
- Mapping all virtual devices to real devices.
- Performing appropriate accounting.
- Informing the user of the date-time of the most recent revision to the system log message (LOGMSG), and of the presence of any outstanding spooled files in his virtual reader, printer, or punch.
- Sending a ready message to the user with the date-time (and weekday), and a message to the system operator indicating the user has logged on.

If the user has a device address or a named system in his user directory and he has not suppressed its initialization via an option on the LOGON command line, then that device or named system is then loaded (via IPL) at the conclusion of the LOGON process. Otherwise, when the LOGON functions are complete, the user is placed in the console function mode with a read on his terminal, ready for the entry of his first desired command.

Under the latter condition of no automatic IPL, the user can IPL an alternate nucleus by using the STOP option in the IPL command. This option will cause the normal IPL procedure to halt execution, prior to loading the initial PSW, and issue a diagnose code 8 placing the user in CP console function mode. A

hexadecimal character entered in location X'08' will change the nucleus name. A hexadecimal character entered in location X'09' will change the apparent storage size. The BEGIN command allows the IPL procedure to continue.

#### User I/O Reconfiguration

Three commands are available to alter the I/O configuration of a user's virtual machine after he has logged on to the system. Two of the commands are available to the user, while the third is restricted to the system operator, since it affects the status of real devices attached to the system. The ATTACH and DETACH commands are contained in DMKVDB and DEFINE in DMKDEF. Both pageable modules are called by the system command scanner (DMKCFM) after their format and privilege classes have been validated. These commands access the same control-block building subroutines in the module DMKVDS that are used by the LOGON processor DMKLOG.

Attaching a Real Device: The system operator can dedicate a real device of any type to a single user by issuing the ATTACH command. The device attached is available only to the given user, and all I/O requests to it are handled via CCW translation. If the device is a DASD, cylinder relocation does not occur when seek addresses or home addresses are referenced. The I/O Supervisor does not queue operations on the device, nor automatically restart it nor do ordered seek queueing. Nonsharable devices such as tape drives must be attached to a user in order to be accessed by a virtual machine. A user can also have a dedicated card read/punch or printer. However, this is usually not necessary because of the unit record spooling facilities of CP. Unit record input or output on a dedicated (attached) device is not spooled by CP. The unit attached may be given a different virtual address than its real address; however, the user may not already have a virtual device at the attached address. A real device cannot be attached (1) if it is currently dedicated to another user, (2) if it contains mini-disks that are in use by other users, or (3) if it is a system owned volume that is in use for spooling or paging.

Defining a Virtual Device: A system user can DEFINE a new virtual device that does not require the dedication of a corresponding real device. Devices that can be defined are consoles, spooled readers, punches and printers, dialable TP lines, virtual channel-to-channel adapters, pseudo timers, and temporary disks. With DEFINE, the user can change any existing virtual device address whether it corresponds to a shared or dedicated real device or no real device unit.

Temporary disks are dynamically obtained cylinders of DASD storage space. They are available to the user for as long as they are part of his virtual configuration, but the data on them is destroyed after the user detaches the area. For all other purposes, however, they appear to be a standard disk.

Detaching a Virtual Device: A virtual device can be removed from a users configuration prior to logout via the DETACH command. A user can detach any of his own devices, and the system operator can detach a real device from a user. In this case, the user is informed of the operator's action. A real device can be detached only if it is dedicated to a single user or is attached to the system and is not in use when the DETACH is issued.

#### VIRTUAL CONSOLE SIMULATION

DMKVCN receives control from the virtual machine I/O executive, DMKVIO. When control is received, the device is available with no interrupts pending. A console control block, VCONCTL, that is obtained from storage and chained from the virtual device control block, VDEVBLOCK, by DMKLOG is accessed for use during the interpretation of the virtual console I/O sequence. The user's CAW is examined for validity. If it is valid, the TRANS macro is issued to fetch the first user CCW. This CCW is moved to the VCONCTL block for analysis.

The CCW is analyzed to determine if it is a read, a write, a control, a sense, a TIC, or an invalid operation. Based upon the analysis, the appropriate processing routine in DMKVCN is invoked.

The Read Simulation Routine: Obtains a buffer for input data from FREE storage. The location of the

buffer is remembered in the VCONCTL block. The DMKQCNRD routine is called to schedule and perform an actual read to the corresponding real device representing the user's virtual console. If SET LINEDIT ON is specified, the buffer data is edited and translated to EBCDIC. When the read is completed, the data is moved to the specified user address obtained from the address portion of the virtual CCW. If command chaining is specified, processing returns to fetch and analyze the next CCW. If command chaining is not specified, the virtual CSW is constructed in the VDEVBLOK and an interrupt is flagged as sending in the VMBLOK.

The Write Simulation Routine: Obtains a buffer for construction of the output message from free storage. The users data is located from the virtual CCW address in the VCONCTL block and moved to the data buffer. The DMKQCNWT routine is called to write the data in the buffer and provide the necessary length, translation, and format functions. Control is received back at the DMKVCN module upon completion of the writing. At this point, the virtual CCW is re-examined. If command chaining is specified, processing continues to fetch and analyze the next CCW. If command chaining is not specified, the virtual CSW is constructed in the VDEVBLOK and an interrupt is flagged as pending in the VMBLOK.

The Control Simulation Routine: Is used for the NOP and ALARM operations. A NOP operation requires no data transfer or I/O operation. An ALARM operation has no equivalent on low speed teleprocessing equipment; thus, a message indicating the alarm operation is constructed. DMKQCNWT is called to output the constructed message. If the command is chained, processing continues (for NOP or ALARM) to fetch the next CCW and analyze it. If command chaining is not specified and this is not the first CCW, a virtual CSW is constructed in the VDEVBLOK and an interrupt is flagged as pending in the VMBLOK. If this is the first (and only) CCW, then a condition code of 1 is presented with channel end and device end in the virtual CSW.

A Virtual Sense Operation: Is similar to a control operation, because no actual I/O operation is performed. However, there is data transfer. The sense data from the VDEVBLOK is moved to the virtual storage location specified in the virtual CCW address. If the command is chained, processing continues to fetch the

next CCW and analyze it. Otherwise, an interrupt is flagged as pending in the VMBLOK.

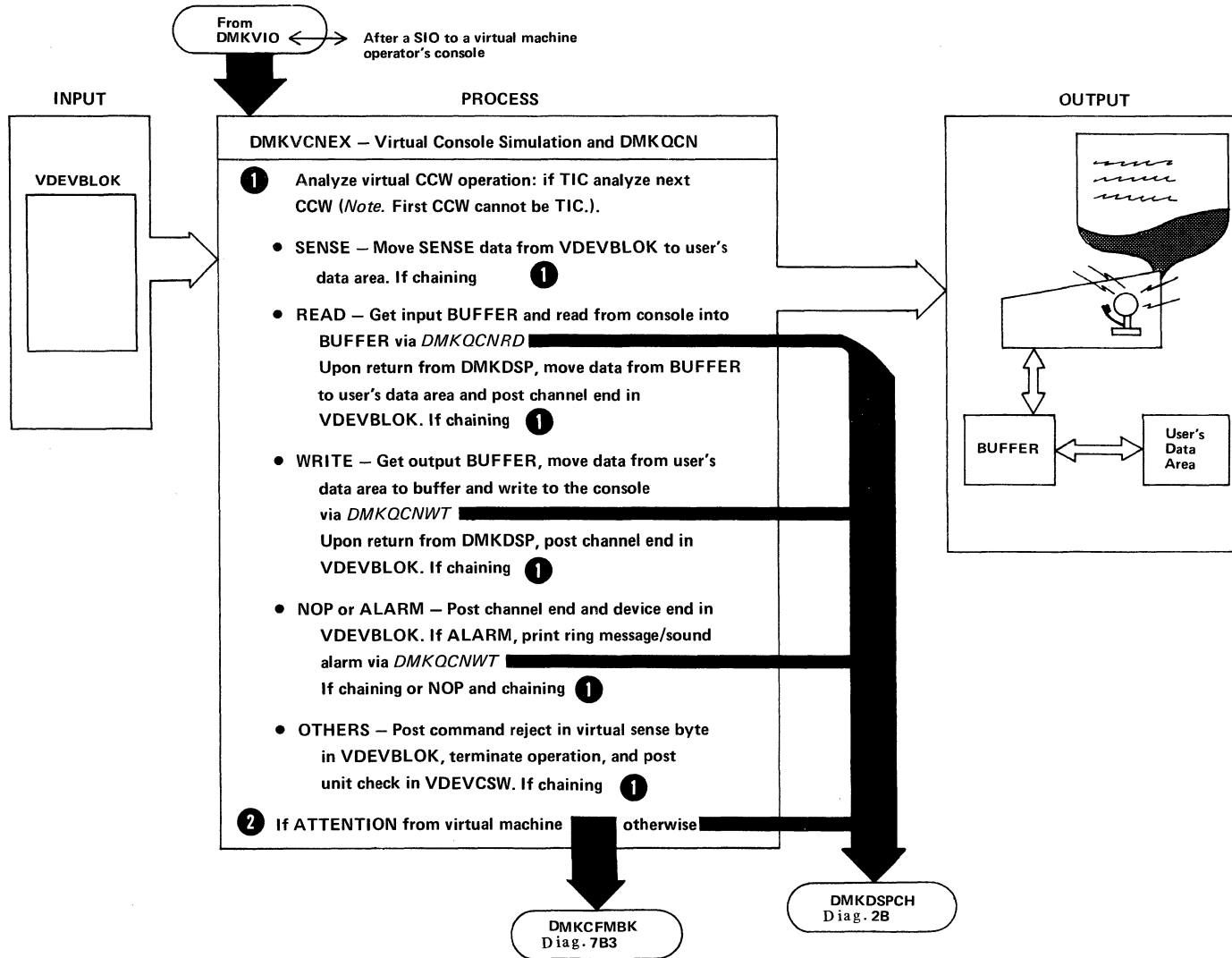
A Virtual TIC: Fetches the virtual CCW addressed by the TIC address and analysis of the fetched CCW continues. If the fetched CCW is itself a TIC, or if the TIC is the first CCW, a channel program check

condition is reflected to the virtual machine as an interrupt or as a CSW stored condition respectively.

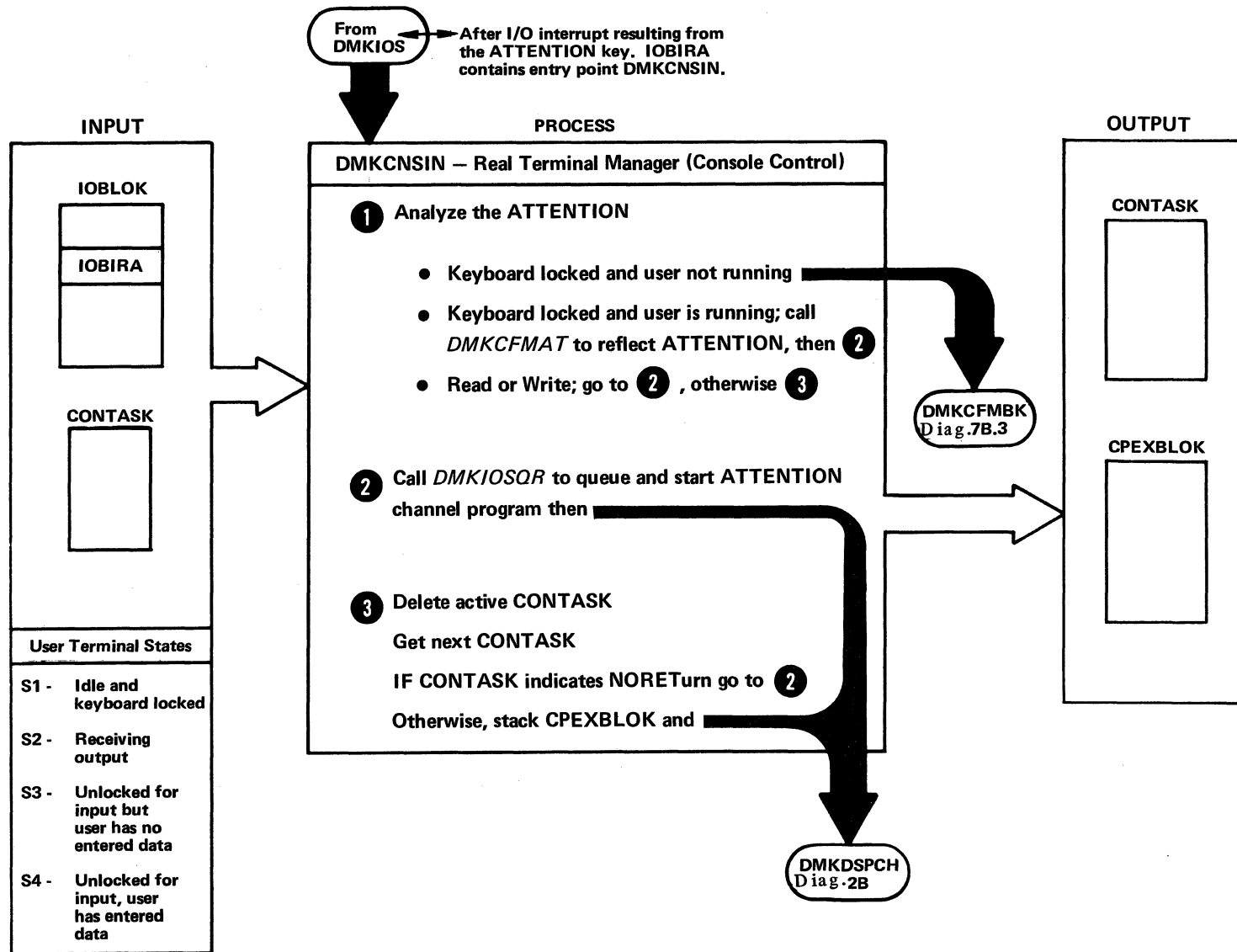
Any other operation is considered invalid. Command reject status is posted in the virtual sense byte and the operation is terminated with unit check status presented in the virtual CSW.



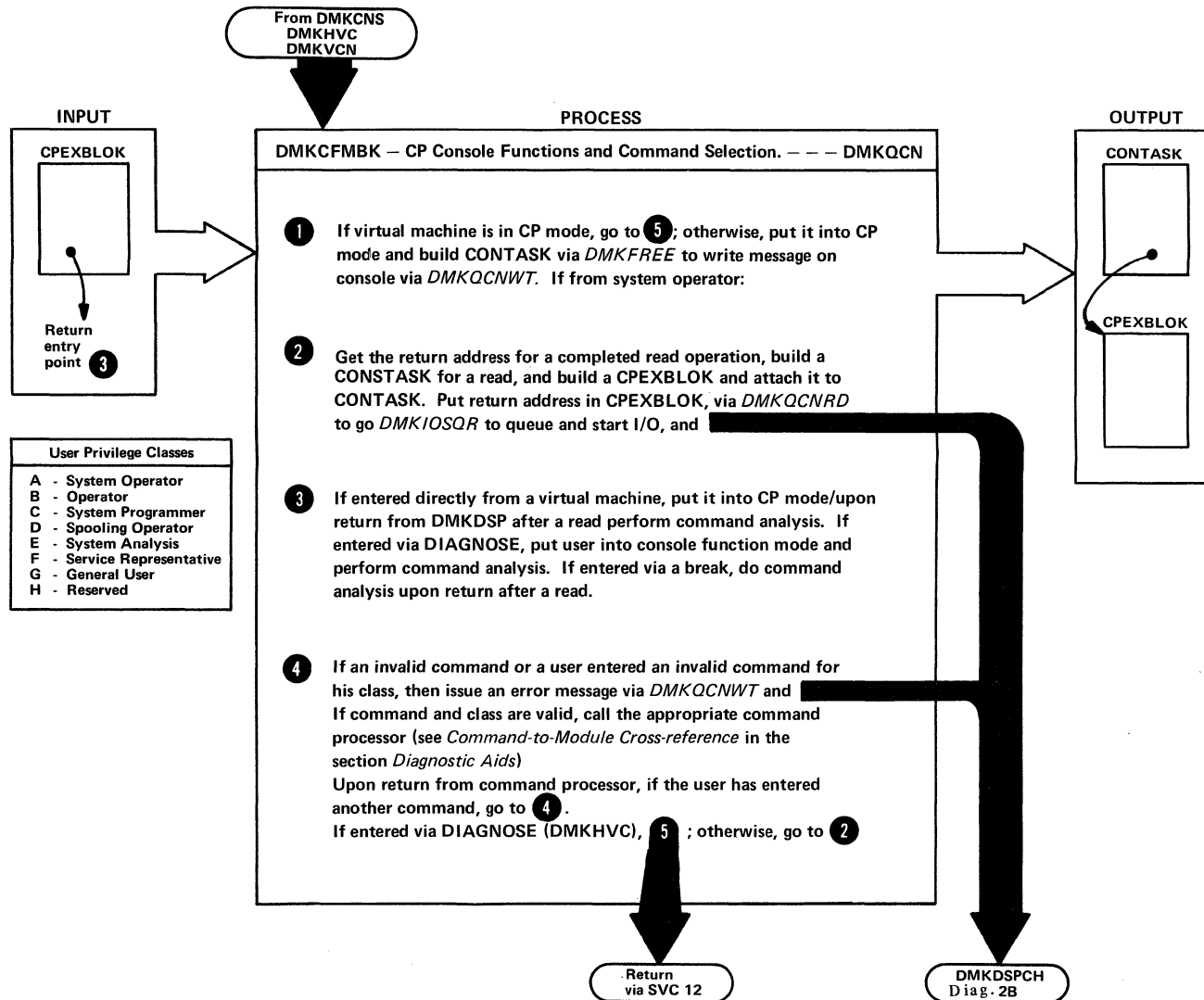
Diag. 7B1. Virtual Console Simulation, Real Terminal Operation



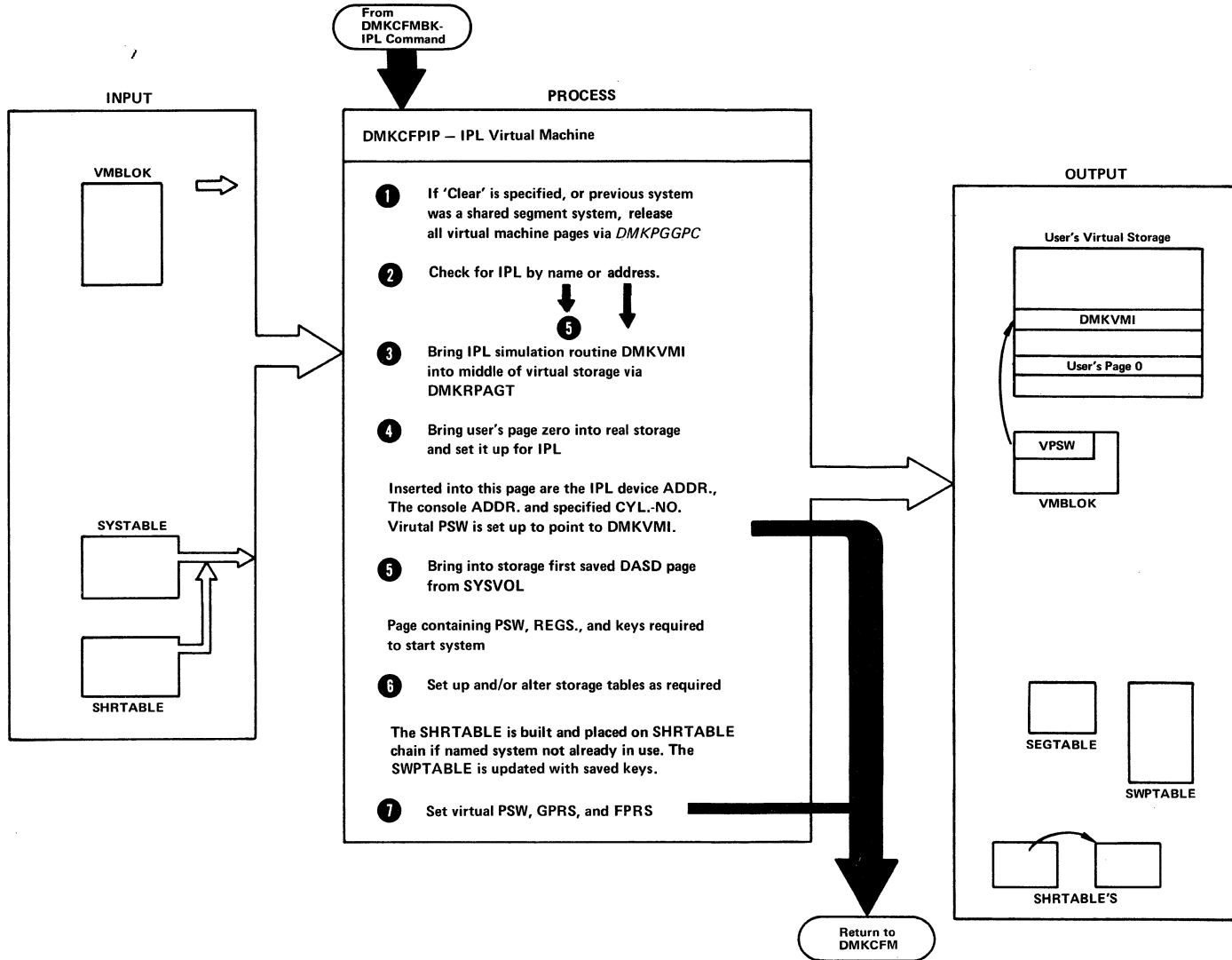
Diag. 7B2. Console Function Control



Diag. 7B3. Function Call Control, Command Selection

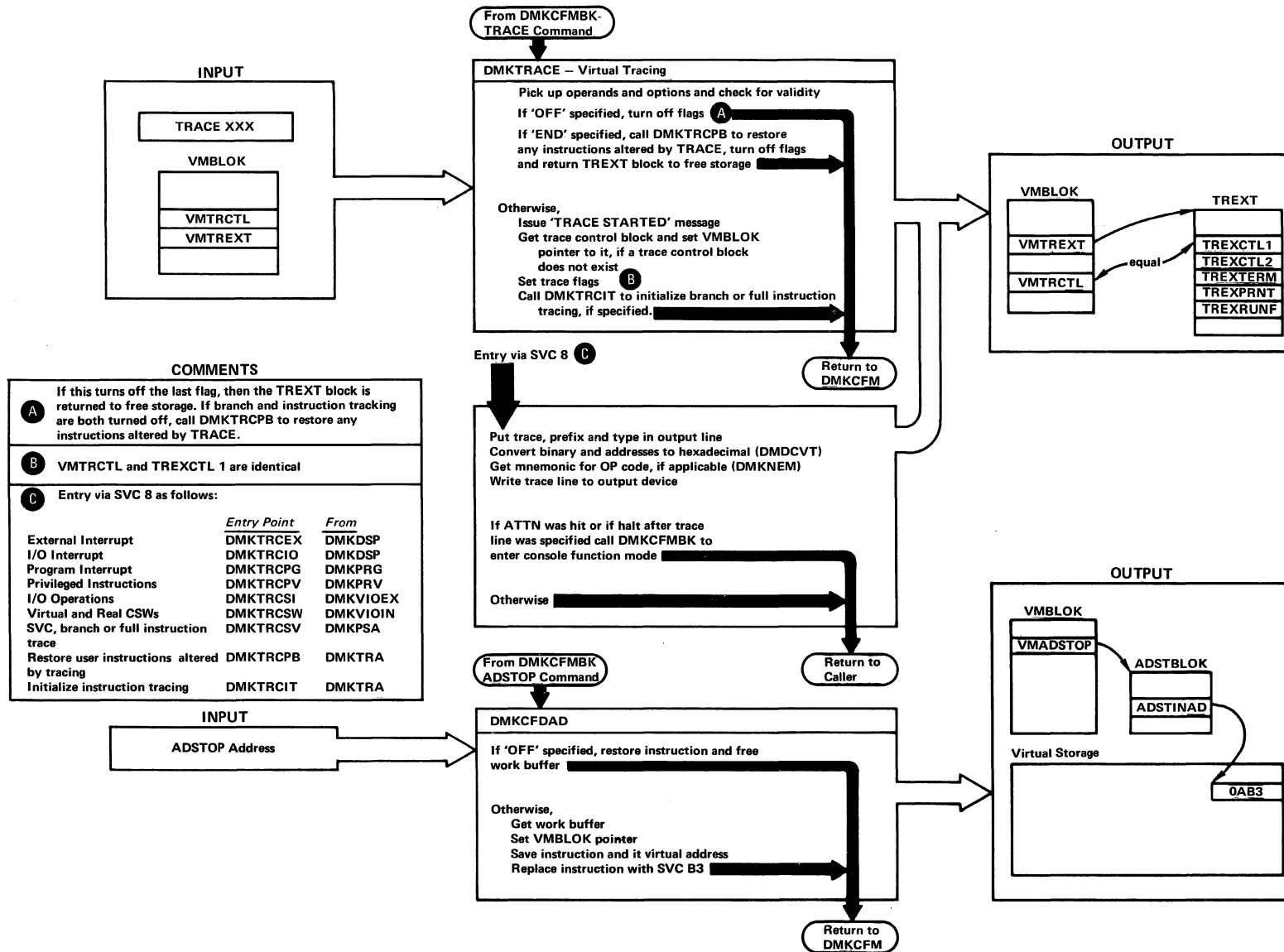


Diag. 7B3.1. Virtual Machine IPL





Diag. 7B3.2. Virtual Tracing



DISCONNECTING A USER: A user may permanently or temporarily disconnect himself from the system by a console command, or he may be forcibly disconnected by the operator or the system. In any case, the routines that handle the termination process are in the pageable module DMKUSO.

Permanent Disconnect: The user may voluntarily exit from the system via the LOGOFF (or LOGOUT) command. This command terminates all virtual machine operation, releases all storage occupied by control blocks and user virtual storage pages, and disconnects the teleprocessing line connection to the user's terminal. If the user specifies the HOLD option with LOGOFF, all of the above occurs, except the teleprocessing line remains enabled. This option is especially useful for dialed connections that will be reused immediately by another user.

The user can be forced off the system by the system operator via the FORCE command. This has the same effect as a user-initiated logoff, except that the user is informed that the operator has logged him off. A user may also be logged off the system:

- If the time for a read of a system password expires (28 seconds).
- If he makes a connection to the system but does not logon within a given period.
- If he is running disconnected (without an active terminal) and his virtual machine attempts a terminal read or enters a disabled wait state.

The LOGOFF command is processed by the DMKUSOLG and DMKUSOFF subroutines. DMKUSOFF is also called directly by DMKDSP to force the logoff of a disconnected user as previously described.

Temporary Disconnect: A user may temporarily disconnect his terminal from his virtual machine while allowing the virtual machine to continue to run via the DISCONN command. This command flags the virtual machine as being disconnected and releases the user's terminal and teleprocessing line. If the HOLD option was specified in the DISCONN command, CP allows the line to remain enabled, and another user can use the terminal to LOGON. The disconnected virtual machine continues to be dispatched until it either attempts to execute a terminal read to the disconnected console or it enters a disabled wait state. At this time, the dispatcher (DMKDSP) calls the routine DMKUSOFF directly to force the machine out of the system. While the machine is disconnected from its virtual console (real terminal) any terminal output is lost; in addition, CP may apply a disconnected penalty to the machines scheduling priority, in order to bias the system in favor of interactive users.

A user may also be disconnected by the system operator. If the disconnected user logs on to the system while his disconnected machine is still running, he is reconnected and can continue to interact with the system in the usual manner.

The DISCONN command is processed by the DMKUSO | subroutine.

## RECOVERY MANAGEMENT SUPPORT (RMS)

The Machine Check Handler (MCH) minimizes the lost computing time due to machine malfunction. MCH does this by attempting to correct the malfunction immediately, and by producing machine check records and messages to assist the service representatives in determining the cause of the problem.

The Channel Check Handler (CCH) aids the Input/Output Supervisor (DMKIOS) in the recovery from channel errors. CCH provides the device dependent Error Recovery Programs (ERPs) with the information needed to retry a channel operation which has failed.

This support is standard and model independent on the external level (from the user's point of view there are no considerations, at system generation time, for model dependencies).

### SYSTEM INITIALIZATION FOR RMS

DMKIOEFL is called by DMKCPI to initialize the error recording at cold start and warm start time. DMKIOEFL will give control to DMKIOG to initialize the MCH area. A store CPU ID (STIDP) instruction is performed to determine if VM/370 is running in a virtual machine environment, or running standalone on the real machine. If VM/370 is running in a virtual machine the version code will be set to a hexadecimal 'FF' by DMKPRV. If the version code returned is hexadecimal 'FF,' the RMS functions will not be initialized beyond putting the wait bit on in the machine check new PSW (virtual). The logic of this is that machine check interrupts and channel errors (other than Channel Data Checks) will not be reflected to any virtual machine. VM/370 running on the real machine will make the determination as to whether the virtual machine should be terminated.

If the version code is not X'FF,' DMKIOG determines what channels are on line by performing a Store Channel ID (STIDC) instruction and saves the channel type for each channel on line. The maximum machine check extended logout length (MCEL) indicated by the Store CPU ID (STIDP) instruction is added to the length of the MCH record header, fixed logout length and damage

assessment data field. DMKIOG will then call DMKFRE to obtain the necessary storage to be allocated for the MCH record area and the CP executing block (CPEXBLOK). DMKIOG saves the pointers for the Machine Check Record and the CPEXBLOK in DMKMCH. DMKIOG obtains the storage for the I/O extended logout area and initializes the logout area and the ECSW to ones. The I/O extended logout pointer is saved at location 172 and control register 15 is initialized with the address of the extended logout area. The length of the CCH record and the online channel types are saved in DMKCCH. It should be noted that the ability of a CPU to produce an extended logout or I/O extended logout and the length of the logouts are both model and channel dependent. If VM/370 is being initialized on a Model 165 II or 168, the 2860, 2870, and 2880 standalone channel modules are loaded and locked by the paging supervisor and the pointers are saved in DMKCCH. If VM/370 is being initialized on any other model, the integrated channel support is assumed; this support is part of the Channel Control Subroutine of DMKCCH. Before returning to DMKIOE the MCH/CCH recording cylinder for error recording is initialized. DMKIOE passes control back to DMKCPI and control register 14 is initialized with the proper mask to record machine checks.

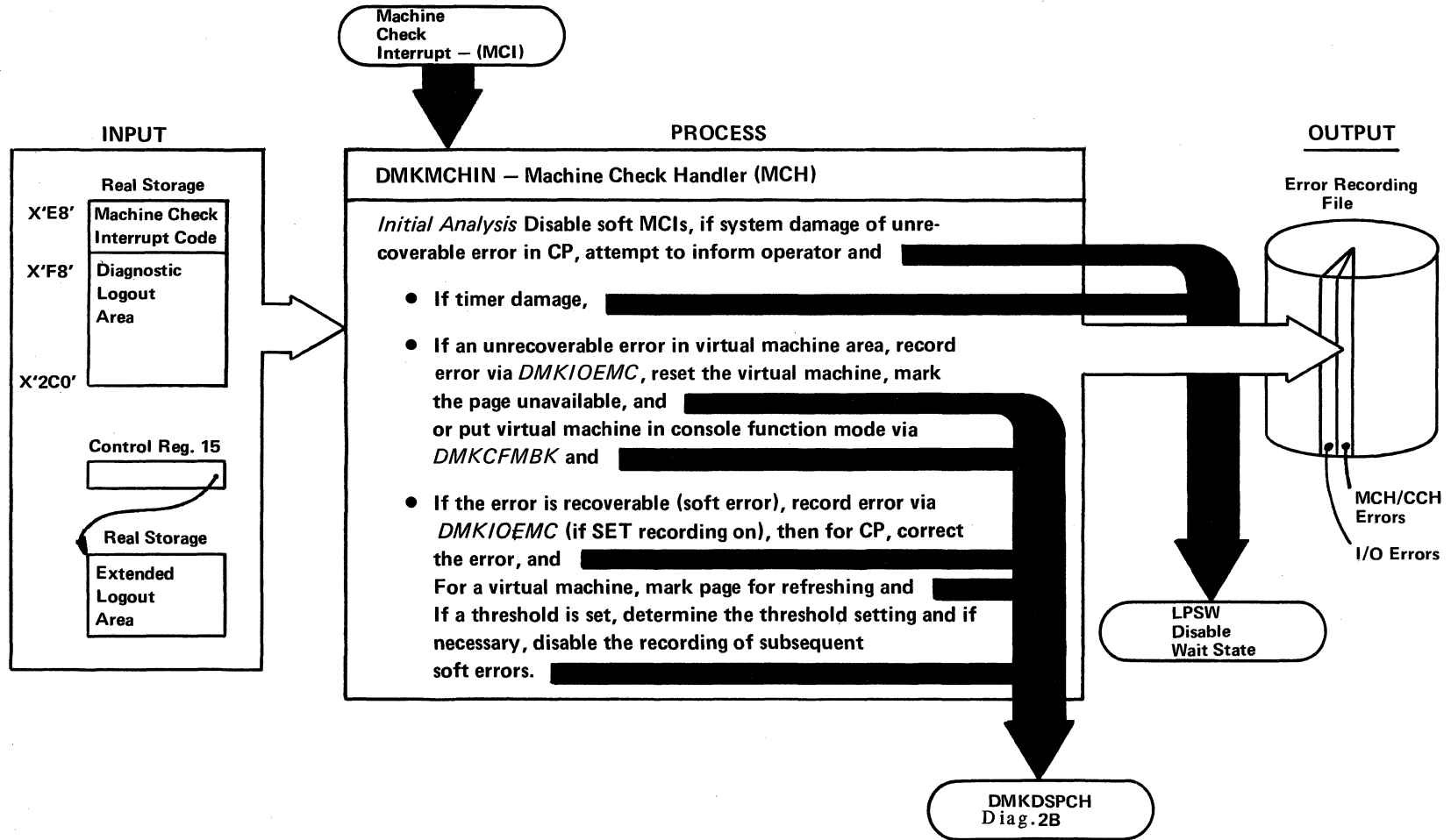
### OVERVIEW OF MACHINE CHECK HANDLER

A machine malfunction can originate from the CPU, real storage or control storage. When any of these fails to work properly, an attempt to correct the malfunction is made by the CPU.

Whenever the malfunction is corrected, the Machine Check Handler (MCH) is notified by a machine check interrupt and the CPU logs out fields of information in real storage, detailing the cause and nature of the error. The model independent data is stored in the fixed logout area and the model dependent data is stored in the extended logout area. The Machine Check Handler uses these fields to analyze the error, format an error record, and write the record out on the error recording cylinder of SYSRES.

If the machine fails to recover from the malfunction through its own recovery facilities, the Machine Check Handler is notified by a machine check interrupt and an

Diag. 8B1. Machine Check Handler (MCH)



interruption code, noting that the recovery attempt was unsuccessful, is inserted in the fixed logout area. The Machine Check Handler then analyzes the data and attempts to keep the system as fully operational as possible.

Recovery from machine malfunctions can be divided into four categories: functional recovery, system recovery, system-supported restart and system repair. These levels of error recovery are discussed in their order of acceptability, functional recovery being most acceptable and system repair being least acceptable:

FUNCTIONAL RECOVERY: Functional recovery is recovery from a machine check without adverse effect on the system or the interrupted user. This type of recovery can be made by the CPU Retry, the ECC facility, or the Machine Check Handler. The CPU Retry and ECC error correcting facilities are discussed separately in this section since they are significant in the total error recovery scheme. Functional recovery by MCH is made by correcting Storage Protect Feature (SPF) Keys and intermittent errors in real storage.

SYSTEM RECOVERY: System recovery is attempted when functional recovery is impossible. System recovery is the continuation of system operations at the expense of the interrupted user, who is terminated. System recovery can only take place if the user in question is not critical to continued system operation. An error in a system routine which is considered to be critical to system operation precludes functional recovery and would require a system-supported restart.

SYSTEM-SUPPORTED RESTART: When the machine check occurs in a critical routine, the primary system operator is notified that the system cannot continue to operate. An automatic reload of the system occurs. This type of recovery is tried when functional and system recovery have failed or could not be tried.

SYSTEM REPAIR: System repair is recovery that requires the services of maintenance personnel and takes place at the discretion of the operator. Usually, the operator has tried to recover by system-supported restart one or more times with no success. An example of this type of error is when a hard error occurs so frequently that system-supported restart is not successful.

## SYSTEM/370 RECOVERY FEATURES

The operation of the Machine Check Handler depends on certain automatic recovery actions taken by the hardware and on logout information given to it by the hardware.

### CPU Retry

CPU errors are automatically retried by microprogram routines. These routines save source data before it is altered by the operation. When the error is detected, a microprogram returns the CPU to the beginning of the operation, or to a point where the operation was executing correctly, and the operation is repeated. After several unsuccessful retries, the error is considered permanent.

### ECC Validity Checking

ECC checks the validity of data from real and control storage, automatically correcting single-bit errors. It also detects multiple-bit errors but does not correct them. Data enters and leaves storage through a storage adapter unit. This unit checks each double word for correct parity in each byte. If a single-bit error is detected, it is corrected. The corrected double word is then sent back into real or control storage and on to the CPU. When a multiple-bit error is detected, a machine check interruption occurs, and the error location is placed in the fixed logout area. MCH gains control and attempts to recover from the error.

Control Registers

Two control registers are used by MCH for loading and storing control information (see Figure 8). Control register 14 contains mask bits which specify whether certain conditions can cause machine check interruptions and mask bits which control conditions under which an extended logout can occur. Control register 15 contains the address of the extended logout area.

Word	Bits	Name of Field	
14	0	Check-Stop Control	Mch-Chk Handling
14	1	Synch. MCEL Ctrl.	Mch-Chk Handling
14	2	I/O Extended Logout Ctrl.	Chan-Chk Handling
14	4	Recovery Report Mask	Mch-Chk Handling
14	5	Degradation Report Mask	Mch-Chk Handling
14	6	External Damage Report Mask	Mch-Chk Handling
14	7	Warning Mask	Mch-Chk Handling
14	8	Asynch. MCEL Control	Mch-Chk Handling
14	9	Asynch. Fixed Log Ctrl.	Mch-Chk Handling
15	8-28	MCEL Address	Mch-Chk Handling

Figure 8. Control Register Assignments for RMS.

Machine Check Handler Subroutines

VM/370 Machine Check Handler (DMKMCH) consists of the following functions:

1. Initial Analysis Subroutine
2. Main Storage Analysis Subroutine
3. SPF Analysis Subroutine
4. Recovery Facility Mode Switching
5. Operator Communication Subroutine
6. Virtual User Termination Subroutine

7. Soft Recording Subroutine
8. Buffer Error Subroutine
9. Term Subroutine

INITIAL ANALYSIS SUBROUTINE: The Initial Analysis Subroutine of DMKMCH receives control via a machine check interruption. To minimize the possibility of losing logout information by recursive machine check interrupts, the machine check new PSW gives control to DMKMCH with the system disabled for further interruptions. There is always a danger that a machine malfunction may occur immediately after DMKMCH is entered and the system is disabled for interruption. Disabling all interruptions is only a temporary measure to give the Initial Analysis Subroutine time to make the following emergency provisions:

1. It disables for soft machine check interruptions. Soft recording will not be enabled until the error is recorded.
2. It saves the contents of the fixed and extended logout areas in the machine check record.
3. It alters the machine check new PSW to point to the Term Subroutine. The Term Subroutine is designed to handle second machine check errors.
4. It enables for hard machine check interrupts.
5. If a virtual user was running when the interrupt occurred, the running status (GPRS, FPRS, PSW, M.C. old PSW, CRS, etc.) is saved in the user's VMBLOK.
6. It initially examines the machine check data for the following types of errors:

```
MCIC=ZERO
PSW invalid
System damage
Timing facilities damage
```

The occurrence of any of these errors is considered uncorrectable by DMKMCH; the primary system operator is informed, the error is

formatted and recorded, and the system is shutdown followed by an automatic restart function.

7. If the instruction processing damage bit is on, it tests for the following types of malfunctions:
  - Multiple-Bit Error in Main Storage -- Control is given to the Main Storage Analysis Subroutine.
  - SPF Key Error -- Control is given to the SPF Analysis Subroutine.
  - Retry failed -- If the CPU was in supervisor state the error is considered uncorrectable and the VM/370 system is terminated. If the CPU was in problem state, the virtual user is reset or terminated and the system continues operation.
8. If the CPU Retry or ECC was successful on a soft error, control is given to the Soft Recording Subroutine to format the record, write it out on the error recording cylinder, and to update the count of soft error occurrences.
9. If external damage was reported, control is given to the Soft Recording Subroutine to format the record and write it out on the error recording cylinder.

MAIN STORAGE ANALYSIS SUBROUTINE: The Main Storage Analysis Subroutine is given control when it is determined that the machine check interrupt was caused by a multiple-bit storage error. An initial function is performed to point the machine check new PSW to an internal subroutine to indicate a solid machine check, in the event of a machine check interrupt while exercising main storage.

Damaged storage areas associated with any portion of the CP nucleus itself cannot be refreshed; multiple-bit storage errors in CP cause the VM/370 system to be terminated. An automatic restart will reinitialize VM/370.

If the damage is not in the CP nucleus, main storage is exercised to determine if the failure is solid or intermittent. If the failure is considered solid, the

4K page frame is marked unavailable for use by the system. If the failure is considered intermittent, the page frame is marked invalid. The change bits associated with the damaged page frame are checked to determine if the page had been altered by the virtual machine. If no alteration had occurred, VM/370 assigns a new page frame to the virtual machine and a backup copy of the page is brought into storage the next time the page is referenced. If the page had been altered VM/370 resets or terminates the virtual machine, clears its virtual storage, and sends an appropriate message to the user. Normal system operation continues for all other users.

STORAGE PROTECT FEATURE (SPF) ANALYSIS SUBROUTINE: The SPF Analysis Subroutine is given control when it is determined that the machine check interrupt was caused by an SPF error. An initial function is performed to point the machine check new PSW to an internal subroutine in the event of a machine check interrupt during testing and validation. The SPF Analysis routine then determines if the error was associated with a failure in the virtual user storage or in the storage associated with CP itself.

An SPF error associated with VM/370 is a potentially catastrophic failure. Namely, VM/370 always runs with a PSW key of zero, which means that the SPF Key in memory is not checked for an out of parity condition. The SPF Analysis Subroutine exercises all sixteen keys in the failing storage 2K page frame. If an SPF machine check occurs in exercising the sixteen keys five times each, then the error is considered solid and the operating system is terminated with a system shutdown. The system is automatically restarted and the VM/370 is reinitialized. If an SPF machine check does not occur, the machine check is considered intermittent. The zero key is restored to the failing storage 2K page frame and this is done transparent to the virtual users.

If an SPF machine check occurs which is associated with a virtual user, the SPF Analysis subroutine exercises all sixteen keys in the failing storage 2K page frame. If an SPF machine check does not occur, then the machine check is considered intermittent and the SWPTABLE for the page associated with the failing storage address is located. The storage key for the failing 2K storage page frame is retrieved from the

SWPTABLE and the change and reference bits are masked on in the storage key. The storage key is then stored into the affected failing storage 2K page frame. If an SPF machine check occurs in exercising the sixteen keys five times each, then the machine check is considered solid and the following actions are taken. (1) The user is selectively reset or terminated by the virtual user termination subroutine. (2) The 4K page frame associated with the failing address is removed as an available system resource. This is accomplished by locating the CORTABLE for the defective page and altering the CORFPNT and CORBPBNT pointers to make the page unavailable to the system. The CORDISA bit in this CORTABLE is set on to identify the reason for the status of this page in a system dump.

RECOVERY FACILITY MODE SWITCHING: The Recovery Facility Mode Switching subroutine (DMKMCHMS) allows the service representative to change the mode that CPU retry and ECC recording are operating in. This subroutine receives control when a user with privilege class 'P' issues some form of the SET MODE command. A check is initially made to determine if this is VM/370 running under VM/370. If this is the case, the request is ignored and control is returned to the calling routine. The format of the MODE command is as follows:

```
SET MODE {RETRY|MAIN} {QUIET|RECORD}
```

RETRY and MAIN imply CPU retry and main storage respectively.

QUIET causes the specified facility to be placed in quiet mode. RECORD causes the count of soft errors to be reset to zero and the specified facility to be placed in record mode.

OPERATOR COMMUNICATION SUBROUTINE: The Operator Communication subroutine is invoked when the integrity of the system has degraded to a point where automatic shutdown and reload of the system has been tried and was unsuccessful, or could not be attempted due to the severity of the hardware failure. A check is first made to determine if the system operator is logged on as a user, next a check is made to determine if the system operator is disconnected. If either of these checks is not affirmative a message cannot be issued directly to the system operator. A LPSW is performed to place the

CPU in a disabled wait state with a recognizable wait state code in the CPU Instruction Counter.

VIRTUAL USER TERMINATION SUBROUTINE: The virtual user termination subroutine is used to selectively reset or terminate a virtual user whose operation has been interrupted by an uncorrectable machine check. First, the user is marked non-dispatchable to prevent the damaged user from running before reset or termination is performed. The machine check record is formatted and DMKIOEMC is called to record the error. Then the user is notified by a call to DMKQCNWT that a machine check has occurred and that his operation will be terminated. The primary system operator is notified of the virtual user termination via a message issued by a call to DMKQCNWT. If the user is running in the virtual=real area, DMKUSO is called to log the user off the system and to return the storage previously allocated to the user and to clear any outstanding user I/O Requests. The hold option of LOGOFF is invoked to allow a user on a dial facility to retain the connection and thus permit LOGON without re-establishing the line connection. However, if the user is running in the virtual area, and DMKCFM is then called to put the virtual user in console function mode, the user must re-initialize the system to commence operation.

SOFT RECORDING SUBROUTINE: The soft recording subroutine performs two basic functions:

1. Format a machine check record and call DMKIOEMC to record the error on the error recording cylinder.
2. Maintain the threshold for CPU RETRY and ECC errors and switch from recording to quiet mode when the threshold value is exceeded. In order to accomplish this, a counter is maintained by DMKMCH for successful CPU retry and corrected ECC events.

CPU Retry Recording Mode: Recording mode (bit 4 of Control Register 14 set to one) is the initialized state, and normal operating state of VM/370 for CPU Retry errors. Recording mode may also be entered by use of the CP SET command. When 12 soft machine checks



have occurred the soft recording subroutine switches the CPU from recording mode to quiet mode. For the purpose of model-independent implementation this is accomplished by setting bit 4 of Control Register 14 to zero. Since in QUIET mode no soft machine check interruptions occur, a switch from quiet mode to recording mode can be made by issuing the SET MODE RETRY|MAIN RECORD command. While in recording mode corrected CPU RETRY|MAIN reports are formatted and recorded on the VM/370 error recording cylinder, but the primary systems operator is not informed of these incidents.

CPU Retry Quiet Mode: Quiet mode (bit 4 of Control Register 14 set to 0) can be entered in one of two ways: (1) when 12 soft machine checks have occurred, or (2) when the SET MODE RETRY QUIET command is executed by a class 'F' user. In this mode, both CPU retry and ECC reporting are disabled. The CPU will remain in quiet mode until the next system IPL (warmstart or cold start) occurs or a SET MODE RETRY|MAIN RECORD command is executed by a class 'F' user.

ECC Reporting Modes: To achieve model independent support, RMS does not set a specific mode for ECC reporting. The mode in which ECC reporting is initialized depends upon the hardware design for each specific CPU model type. For the IBM System/370 Models 135, 145, 158, and 168 the hardware initialized state (therefore the normal operational state for VM/370) is QUIET mode. For the IBM System/370 Models 155 II, and 165 II the hardware initialized state (therefore the normal operational state for VM/370) is RECORD mode. An automatic restart incident due to a VM/370 failure does not RESET the ECC reporting mode in effect at the time of failure.

The change from RECORD to QUIET mode for ECC reporting can be initiated in either of the following ways; (1) by issuing the SET MODE {MAIN|RETRY} QUIET command, or (2) automatically whenever 12 soft machine checks have

occurred. For the purpose of model independent implementation this will be accomplished by setting bit 4 of Control Register 14 to zero.

The change from QUIET to RECORD mode for ECC reporting can be accomplished by use of the SET MODE MAIN RECORD command. This recording mode option is for use by maintenance personnel only. It should be noted that CPU RETRY is placed in recording mode if it is not in that state when the SET MODE MAIN RECORD command is issued.

While in RECORDING mode, corrected ECC reports are formatted and recorded on the error recording cylinder, but the primary systems operator is not informed of these incidents.

BUFFER ERROR SUBROUTINE: On CPU models equipped with a high speed buffer (155 II, 158, 165 II, 168) or a Data Look Aside Table (DLAT) (165 II, 168) the deletion of buffer blocks due to hardware failure is reported via a DEGRADATION REPORT machine check interrupt. MCH enables itself for degradation report machine check interrupts at system initialization by setting bit 5 of Control Register 14 to 1. If a machine check interrupt occurs which indicates high speed buffer or DLAT damage, MCH formats the record and calls DMKIOEMC to record it on the error recording cylinder, informs the primary systems operator of the failure, and returns control to the system to continue normal operation.

TERM SUBROUTINE: The Term Subroutine is given control in the event of a hard machine check interrupt while DMKMCH is in the process of handling a machine check interrupt. Note that soft error reporting is disabled for the entire time that MCH is processing an error.

An analysis is performed of the machine check interrupt code of the first error to determine if it was a soft error, and if it was, the first error is recorded, the system status is restored and control is restored to the point where the first error occurred. If the first error was a hard error, the Operator Communication Subroutine is given control to issue a message directly to the system operator, and to terminate CP operation.

## OVERVIEW OF CHANNEL CHECK HANDLER

The Channel Check Handler (CCH) aids the I/O Supervisor in recovering from channel errors and informs the operator or service representative of the occurrence of channel errors.

CCH receives control from the I/O Supervisor when a channel data check, channel control check, or interface control check occurs. CCH produces an I/O Error Block (IOERBLOK) for the error recovery program and a record to be written on the error recording cylinder for the system operator or service representative. The operator or service representative may obtain a copy of the record by using the CPREP programs. A message about the channel error is issued each time a record is written on the error recording cylinder.

When the Input/Output Supervisor program detects a channel error during routine status examination following an SIO, TIO, HIO, or an I/O interruption it passes control to the Channel Check Handler (DMKCCH). DMKCCH analyzes the channel logout information and constructs an IOERBLOK, if the error is a channel control or interface control check, and an ECSW will be constructed and placed in the IOERBLOK. The IOERBLOK provides information for the device dependent error recovery procedures. DMKCCH also constructs a record to be recorded on the error recording cylinder. Normally, DMKCCH returns control to the I/O Supervisor after constructing an IOERBLOK and a record. However, if DMKCCH determines that system integrity has been damaged (system reset or invalid unit address, etc.) then CP operation will be terminated. The action taken by DMKCCH for CP termination will be to issue a message directly to the system operator and place the CPU in a disabled wait state with a recognizable wait code in the CPU instruction counter.

Recovery will not be initiated for channel errors associated with I/O events initiated by a virtual user, however these will cause termination of the user after he has been notified of the failure. The error will be recorded by DMKIOECC on the error recording cylinder.

Normally, when DMKCCH returns control to the I/O supervisor, the error recovery program for the device which experienced the error is scheduled. When the ERP

receives control, it prepares to retry the operation if analysis of the IOERBLOK indicates that retry is possible. Depending on the device type and error condition, the ERP will either effect recovery or mark the event fatal and return control to the I/O Supervisor. The I/O Supervisor will call the recording routine DMKIOE to record the channel error.

The primary system operator will be notified of the failure, and DMKIOE will return control to the system and normal processing will continue.

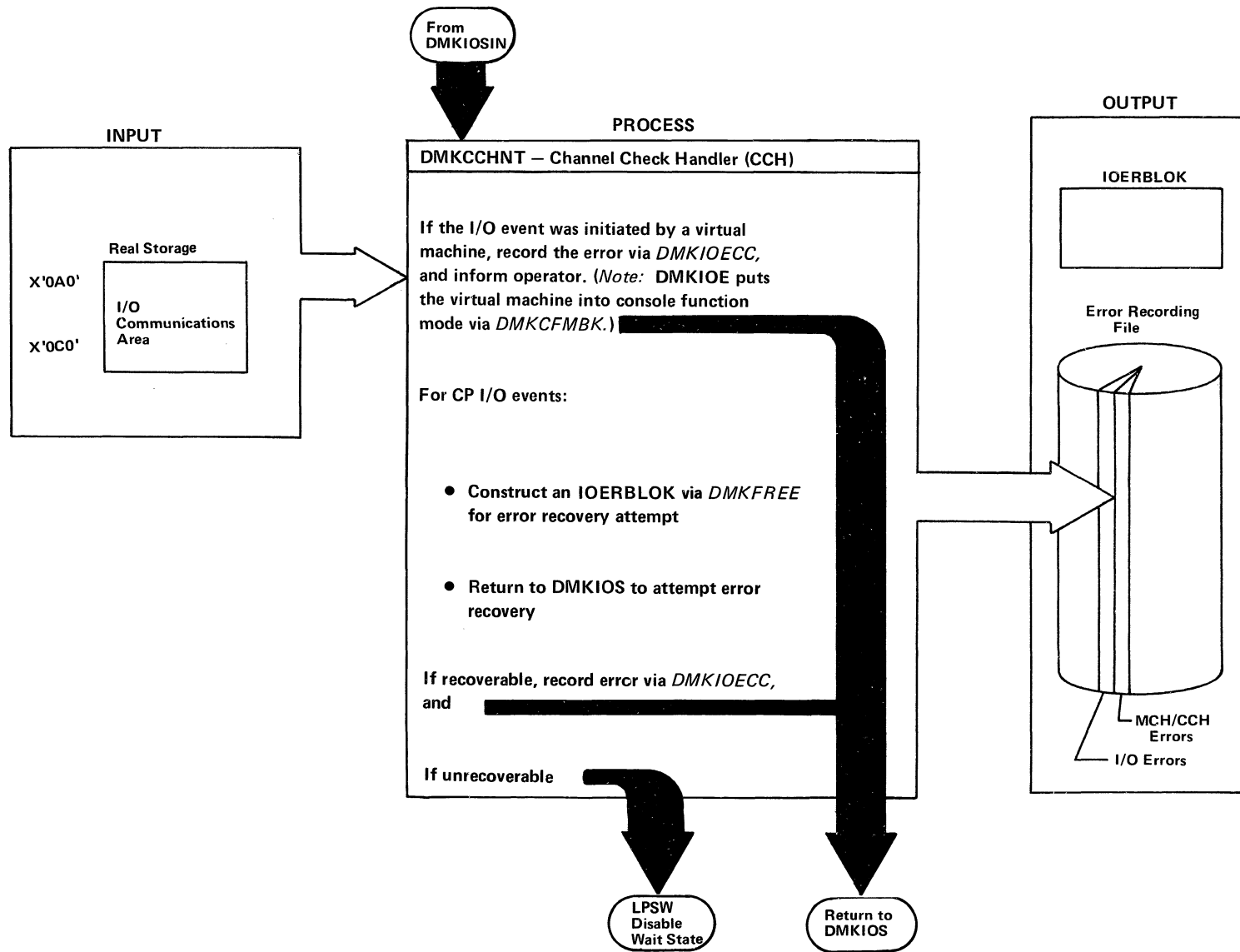
Channel Control Subroutine

Control is passed to the Channel Control Subroutine of DMKCCH after a SIO with failing status stored, or an I/O interrupt due to a channel control check, interface control check, or channel data check.

If "logout pending" is indicated in the CSW, the CP termination flag is set. The existence of real device blocks (RCHBLOK, RCUBLOK, RDEVBLOK), for the failing device address, is determined by a call to DMKSCNRU and an indicator is set if they do exist. An indicator is also set if the IOBLOK for the failing device address exists. A call to DMKFREE obtains storage space for the channel check record and the channel control subroutine builds the record. If the indicators show that the real device blocks and the IOBLOK exist, a call to DMKFREE obtains storage space and the channel control subroutine builds the I/O error block (IOERBLOK); if these blocks do not exist, the IOERBLOK is not built. The IOERBLOK is used for two purposes:

1. The device dependent Error Recording Program (ERP) uses the IOERBLOK to attempt recovery on CP initiated I/O events. If the I/O events that resulted in a channel check are associated with a virtual user, the I/O fatal flag is set in the IOBLOK and the user's virtual machine is reset, cleared, and put into console function mode with a read up on the line. The length and address of the channel check record is placed in the IOERBLOK and the IOERBLOK is chained off the IOBLOK.
2. DMKIOECC uses the IOERBLOK to record the channel check record on the error recording cylinder.

Diag. 8B2. Channel Check Handler (CCH)



The channel control subroutine gives control to a channel dependent error analysis routine to build or save the extended channel status word (ECSW). When the Channel Control Subroutine regains control, eight active addresses are saved in the channel check record.

If the CP termination flag is set, the I/O extended logout data from the channel check record is restored to main storage for use by SEREP. If the system operator is both logged on as a user and connected to the system, a message (DMKCCH603W) is sent to him advising him of the channel error. A LPSW is then executed to place the CPU in a disabled wait state with a wait state code of 002 in the CPU instruction counter.

If the CP termination flag is not set, a check is made to determine if an IOERBLOK was built by the channel control subroutine.

If an IOERBLOK was not built, DMKIOECC is called to record the channel check record on the error recording cylinder. The system operator is then sent a message (DMKCCH601I or DMKCCH602I) informing him of the error and control is then returned to DMKIOS to continue system operation.

If an IOERBLOK was built, control is returned to DMKIOS which calls the appropriate ERP. Whether or not recovery is successful, DMKIOS eventually calls DMKIOE to record the channel check record. DMKIOE examines the status of the error CSW in the IOERBLOK to determine if it was a channel error; if so, it finds the length and pointer to the channel check record and records the error on the error recording cylinder. If this was not a channel error, DMKIOE continues normal processing.

#### Individual Routines

A separate channel error analysis routine is provided for each type of channel for which DMKCCH can be used. The purpose of these routines and the Channel Control Subroutine is to analyze the channel logout to determine the extent of damage and to create a sequence and termination code to be placed in the ECSW in the

IOERBLOK. At system initialization time the correct model dependent channel recovery routine is loaded and the storage necessary to support the routine is allocated. The model dependent error analysis subroutines and routines and their functions are as follows:

INTEGRATED CHANNELS (Models 135, 145, 155 II, 158): Since all of these systems have integrated channels one common subroutine is used to handle all of these CPU types. This subroutine:

- Indicates CP termination if the ECSW is not complete, the channel has been reset, or reset codes are invalid
- Moves the ECSW to the IOERBLOK
- Moves the hardware stored unit address and the I/O extended logout to the channel check record
- Sets the I/O extended logout area and ECSW area to ones
- Returns control to the Channel Control Subroutine

2860 CHANNEL (Models 165 II, 168): The 2860 logout area is checked to determine if a complete logout exists; if not, CP termination is necessary.

A check is made in the logout area for validity of the CSW fields and bits are set in the channel check record's ECSW field to indicate bad fields.

The channel logout is then checked and sequence codes are set based on the presence of a channel control check, or an interface control check. If a channel control check is present, the codes set are determined through parity. The count determines if parity is good and sets a resultant condition code.

The logout area is examined to ensure that the unit address has valid parity and is the same address passed by DMKIOS. If so, the "unit address valid" bit in the ECSW is set. If the unit address is not valid the "unit address valid" bit is reset to indicate the invalid condition.

The ECSW field in the channel check record is moved to the IOERBLOK, if one exists.

After completing the ECSW the 2680 routine moves the 2860 I/O extended logout into the channel check record,

set the I/O extended logout area to ones, and returns to the Channel Control Subroutine.

2870 CHANNEL (Models 165 II, 168): If the channel failed to logout completely, at least part of the logout area is all ones. If a full word of ones is found, a CP termination condition exists.

A check is made in the logout area for valid CSW fields, and bits are set in the channel check record's ECSW field to indicate bad fields.

The termination and sequence codes are set depending on the presence of an interface control check or channel control check. If a channel control check is present, the codes set are determined through parity, count, and/or data transfer checks. For the 2870, parity can be determined directly from the channel logout.

The logout area is also examined to ensure valid parity in the unit address and to ensure that the address is the same as that passed to DMKCCH by DMKIOS. If so, the "unit address valid" bit in the ECSW is set.

The 3rd word of the logout area is also analyzed for type II errors. If one of these type II errors is found, a CP termination condition exists.

The ECSW field in the channel check record is moved to the IOERBLOK, if one exists.

Before returning to the Channel Control Subroutine, the 2870 routine moves the 2870 I/O extended logout into the channel check record and sets the I/O extended logout area to ones.

2880 CHANNEL (Models 165 II and 168): This routine will analyze 9 words of the 28 word logout.

The 2880 Analysis routine handles channel data checks, interface control checks, and channel control checks.

Termination code 3 (system reset) is not set in the ECSW because the 2880 channel does not issue system reset to the devices. Retry codes of zero to five are possible.

Note: There are several catastrophic conditions under which the CP termination flag can be set, in the 2880 analysis routine. They are:

1. The channel did not complete the logout.
2. The CSW is not reliable.
3. The unit address in the I/O interrupt device address field is not correct.

Only a channel check record is needed if the channel has recognized an internal error and has recovered from it without any damage. No recovery action is necessary in these cases.

If the channel address in the I/O interrupt device address field does not match the channel address in the logout, a CP termination condition exists.

If the channel was doing a scan and the unit control word had a parity check a CP termination condition exists. If there was no parity check, there was no damage during the scan and only a channel check record is required.

Depending on the sequence the channel has entered, the termination and sequence codes are set; command address, unit address, and unit status validity is determined; and the sequence code is set valid. The ECSW field in the channel check record is moved into the IOERBLOK, if one exists.

Before returning to the Channel Control Subroutine, the 2880 routine will move the I/O extended logout into the channel check record and set the I/O extended logout area to ones.

#### ERROR RECORDING AND RECOVERY

The error recording facility is made up of three modules. One module (DMKIOE) is resident and the other two (DMKIOF and DMKIOG) are pageable.

The error recording routines records: unit checks, machine checks, channel checks, and hardware environmental counter sense data on the error recording cylinders of the system resident device in a format

suitable for subsequent processing by the CPEREP program. The recorder also initializes the error recording cylinders at IPL time if they are in an unrecognizable format.

When the recorder is entered from DMKIOS, it is entered at DMKIOERR. This entry is used for unit checks and channel data checks. A test is made of the failing CSW (located in the IOERBLOK) to see if the error was a channel error. If it was, control is passed to routine for recording channel checks.

The IOERBLOK sense data, IOBLOK flags, and VMBLOK user class are examined to determine if the error should be recorded. See the section "Errors Recorded" for those that are recorded.

#### Writing the Record

After an error record is formatted, it is added to the error recording cylinder using DMKRPAGT and DMKRPAPT. The error recording cylinders have page sized records (4096 bytes). Each page contains a header (8 bytes) which signifies cylinder and page number of the page (4 bytes), next available space for recording within page (2 bytes), a page in-use indicator (1 byte), and a flag byte. Each record within the page is recorded with a 4-byte length prefix.

If an error record is too large to be added into a page, a new page is retrieved, updated with record, and placed back on the error recording cylinder with the

paging routines.

Two cylinders are used for error recording: one cylinder is used exclusively for recording the I/O errors and the other cylinder for recording MCH/CCH errors. The cylinders that are used for error recording are specified by the user at system generation time. If either error recording cylinder becomes 90 per cent full, a message is issued to the operator using DMKQCNWT to warn him of the condition. If either cylinder becomes full, another message is issued to inform the operator and recording is stopped on that cylinder. Recording continues on the cylinder that is not full.

If a channel check error is to be recorded, the recorder is entered at DMKIOERR or DMKIOECC. The channel check handler determines the entry. A channel check error record is formatted.

A machine check enters at DMKIOEMC. Pointers are passed from the machine check handler in registers 6 and 7 to locate a buffer where the machine check record and length are saved. A machine check error record is recorded with the saved machine check logout and additional information. The machine check error record is written onto the error recording cylinder by using the paging routines.

Hardware environmental counter records are formed using routine DMKIOEEV. This routine is scheduled by DMKIOS after control is returned from the ERP. Sense data information is stored in the IOERBLOK by the ERP. The record formed is called a nonstandard record.

### Errors Recorded

In addition to recording environmental data, the following types of errors are recorded for DASD virtual machines that are not Class F.

- Bus-out check
- Overrun check
- Seek check
- Track overrun check
- Missing address marker check
- Equipment check
- Permanent data check
- Unrecoverable error for a control program initiated channel program.

The following 3420 errors are recorded if the user is not Class F:

- Bus-out check
- Overrun check
- C compare check
- Write TRIG VRC check

- Feed thru check
- Vel/RESTART check
- Velocity change check
- Equipment check

CLEAR and FORMAT Recording Area: DMKIOEFM is called by the CPEREP program via a DIAGNOSE instruction. DMKIOEFM is invoked to reset the specified error recording cylinders (if CLEARALL, CLEARIO, or CLEARMC was specified). The clear is performed by resetting each page-header space-available field. A pointer in storage is then updated to point to the first page on the error recording cylinder available for recording MCH and CCH records and the first page available on the other error recording cylinder for recording outboard errors. Control is then returned to the calling routine.

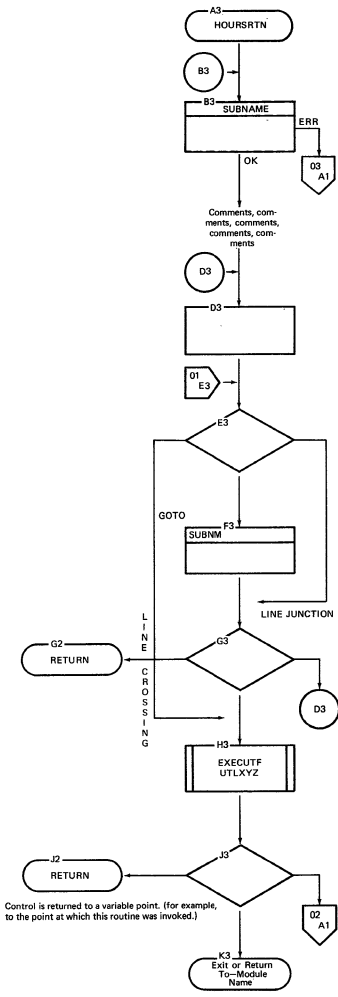
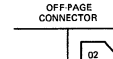
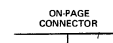
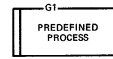
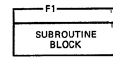
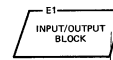
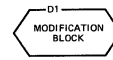
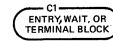
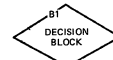
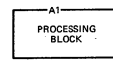
Finding First Recording Cylinder at IPL Time: DMKIOEFL is called by DMKCPI to find the first available page that can be used for error recording. The paging routines, DMKRPAPT and DMKRPAGT, are used to read the error recording cylinder's pages (4096 byte records). As each page record is read it is examined to see if this record is the last recorded. If so, a pointer in storage is saved so recording can continue on that page record. Control is then returned to the caller. If either error recording cylinder is in an unrecognizable format, that cylinder is automatically reformatted by CP.

**PROGRAM ORGANIZATION**

This section contains the flowcharts for all processing modules. The modules are in alphabetical order. To determine the pertinent information about a module, see the Directory entries DMKACO to DMKWRM.

**FLOWCHARTS**

FUNCTIONAL SYMBOLS



The terminal block is used to show entry and exit points of a routine or subroutine. Block A3 shows an entry point named HOURSRTN.

This instruction at this location calls a subroutine called SUBNAME. Upon return from the subroutine, if an error was detected (ERR), then processing resumes at the block A1, part 3 of this series of charts. If an error was not detected (OK), then normal processing resumes.

Comments are included to provide additional information about the logic being displayed.

On-page entry connector one or more branches to this block appear on this part of the flowchart.

Off-page entry connector. A branch to this block appears on another part(s) of this flowchart.

The instruction at location GOTO calls either a subroutine within this module or another module that is used like a subroutine. If the name in the subroutine block begins with DMK, then the call is to another module (see the "Module/Entry Point Directory" to determine the flowchart location of the called module). If the name does not begin with DMK, then the subroutine is in the flowchart series for this module (see the "Subroutine Directory" to determine the flowchart-part location of the called subroutine).

One-page exit connector control branches to block D3 on this part of the flowchart.

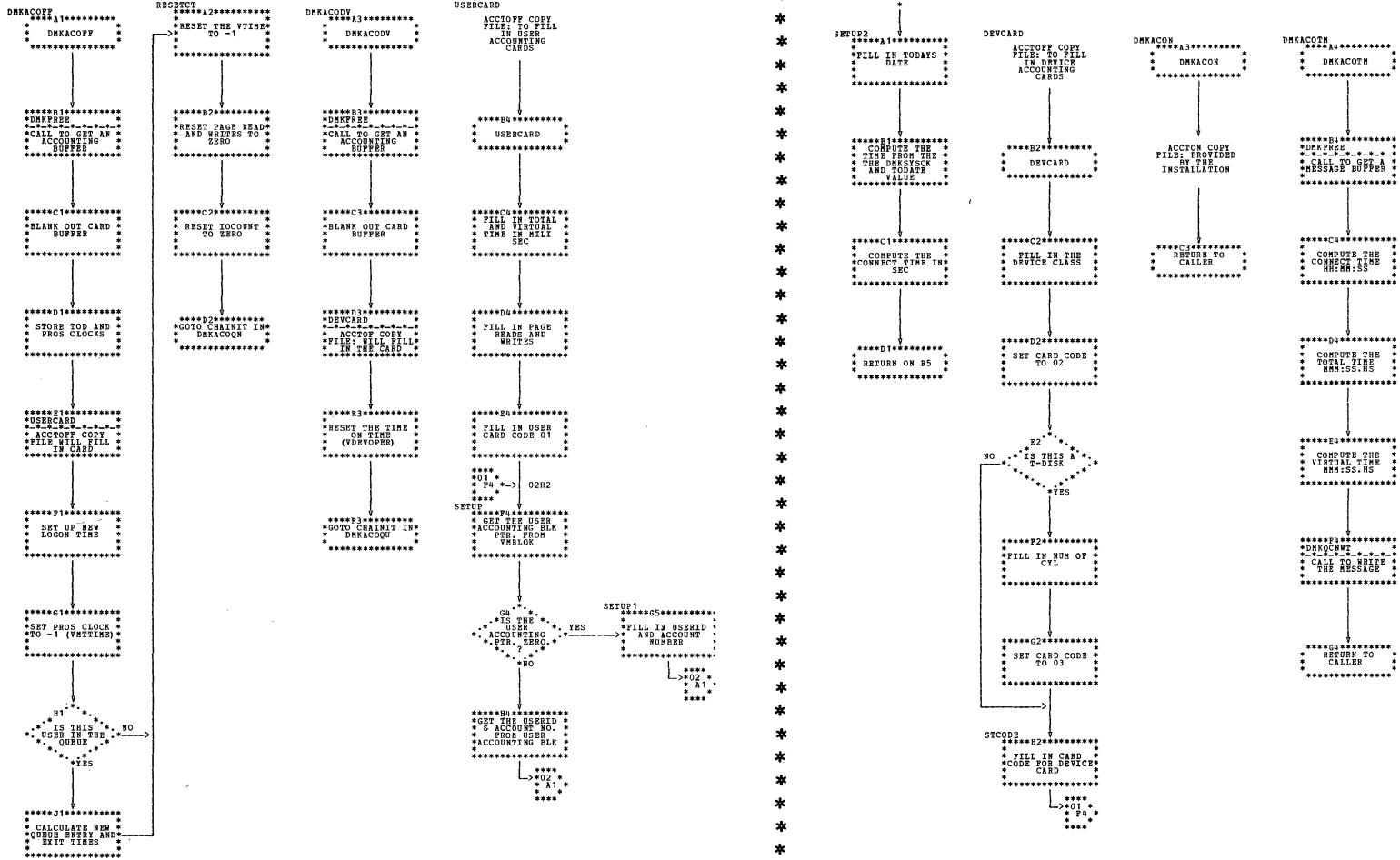
This block refers to a routine or program that is documented in some other publication.

Off-page exit connector control branches to block A1 on part 2 of this series of flowchart.

Control is returned to a variable point. (For example, to the point at which this routine was invoked.)

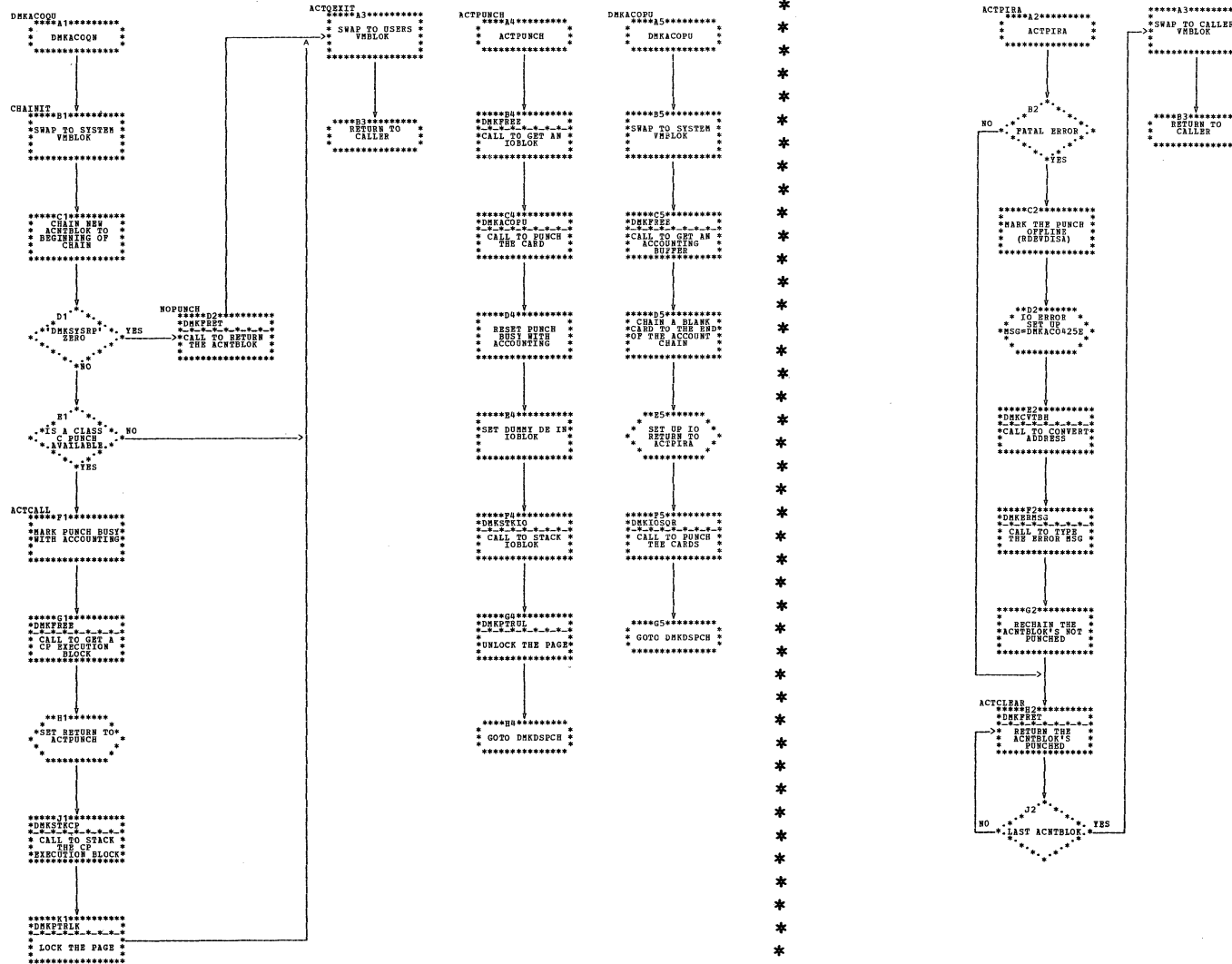
Control branches to an entry point on another flowchart.





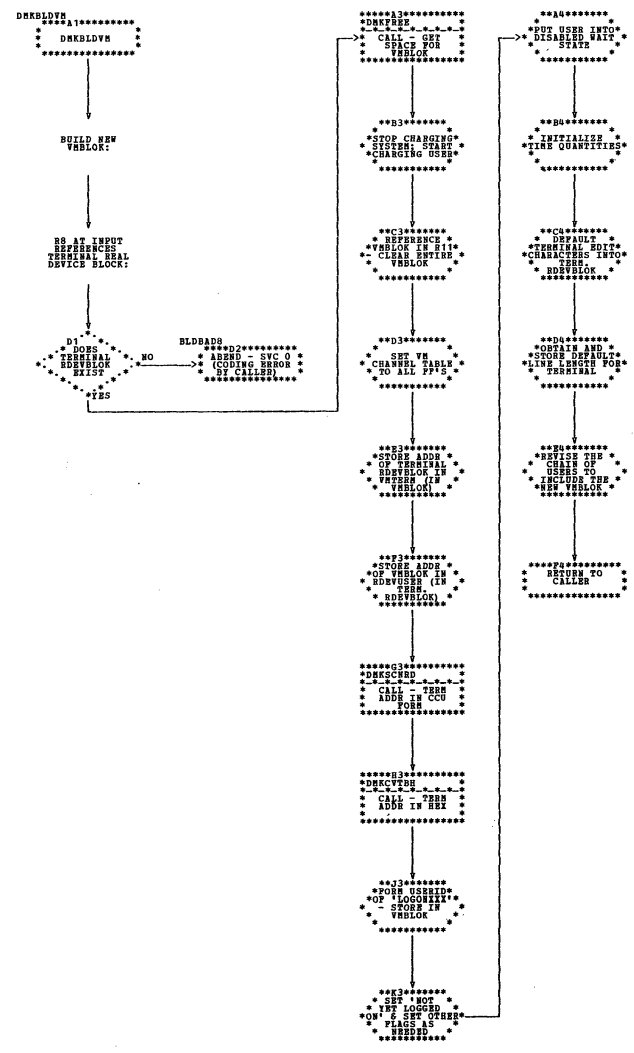
| DMKACO -- Accounting Routines (Parts 1 and 2 of 4)

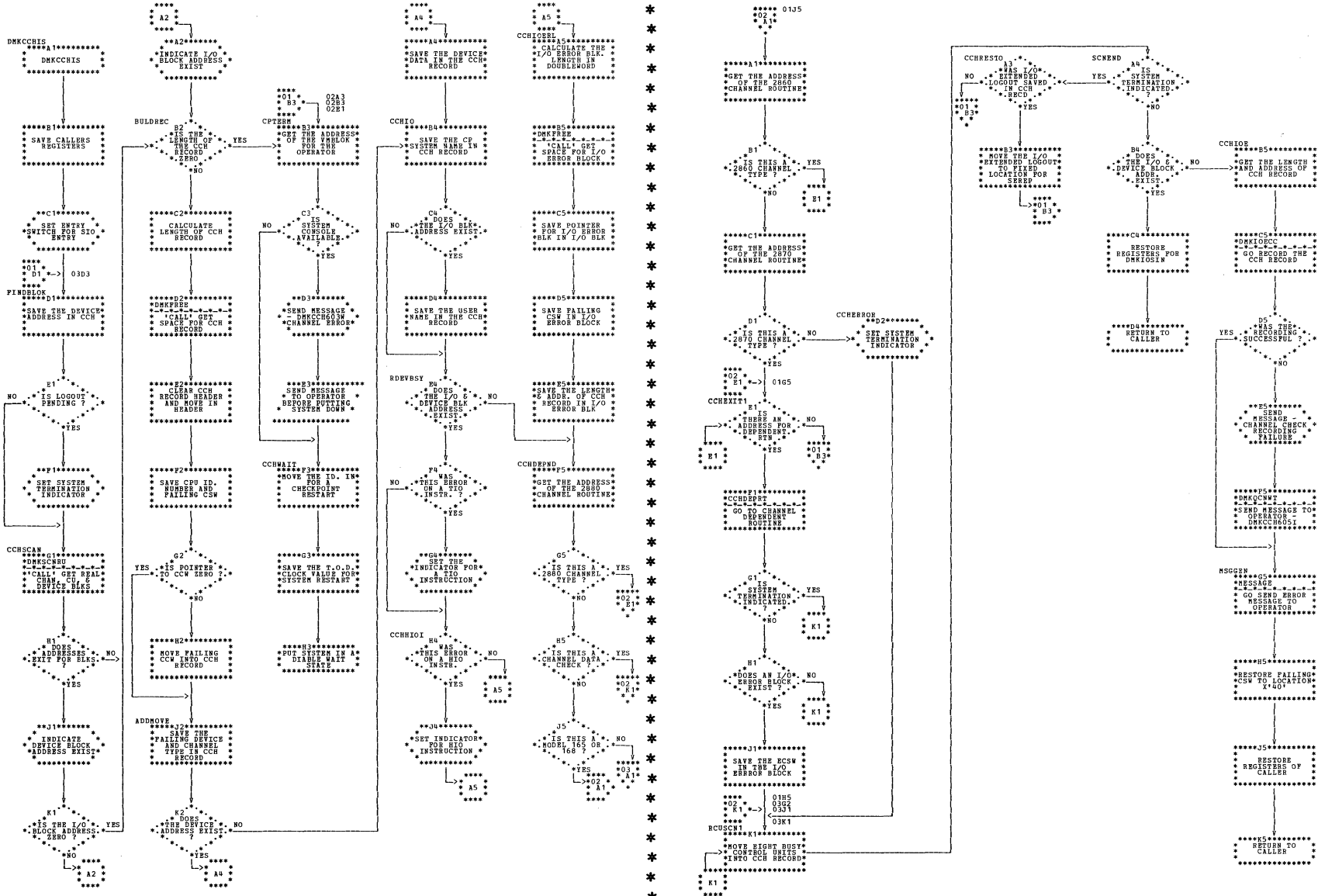
| DMKACCO -- Accounting Routines (Parts 3 and 4 of 4)





DMKELD -- Build/Release Real Storage Tables; Build VMBLOK (Part 3 of 3)





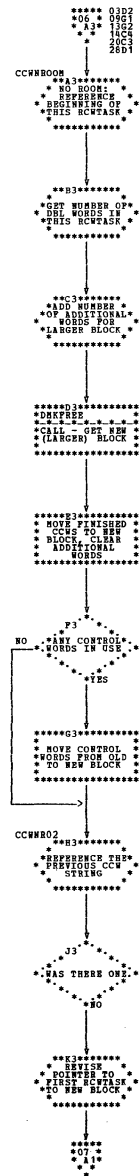
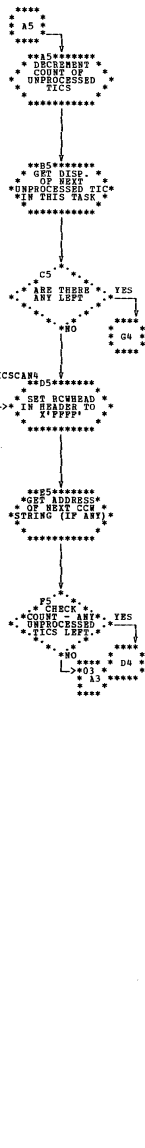
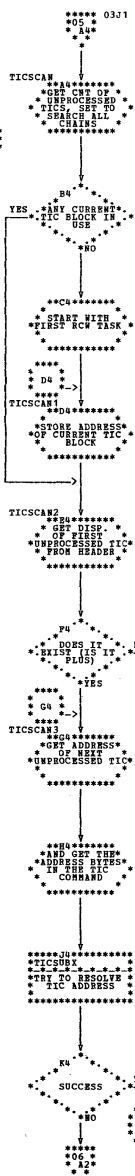
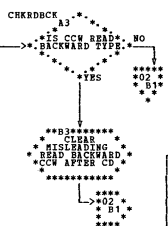
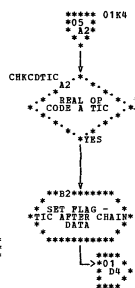
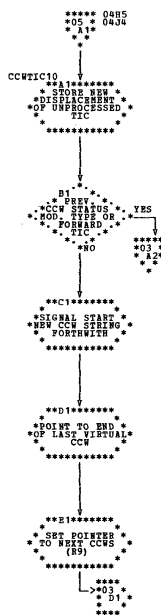
DMKCCCH -- Channel Check Handler (Parts 1 and 2 of 4)





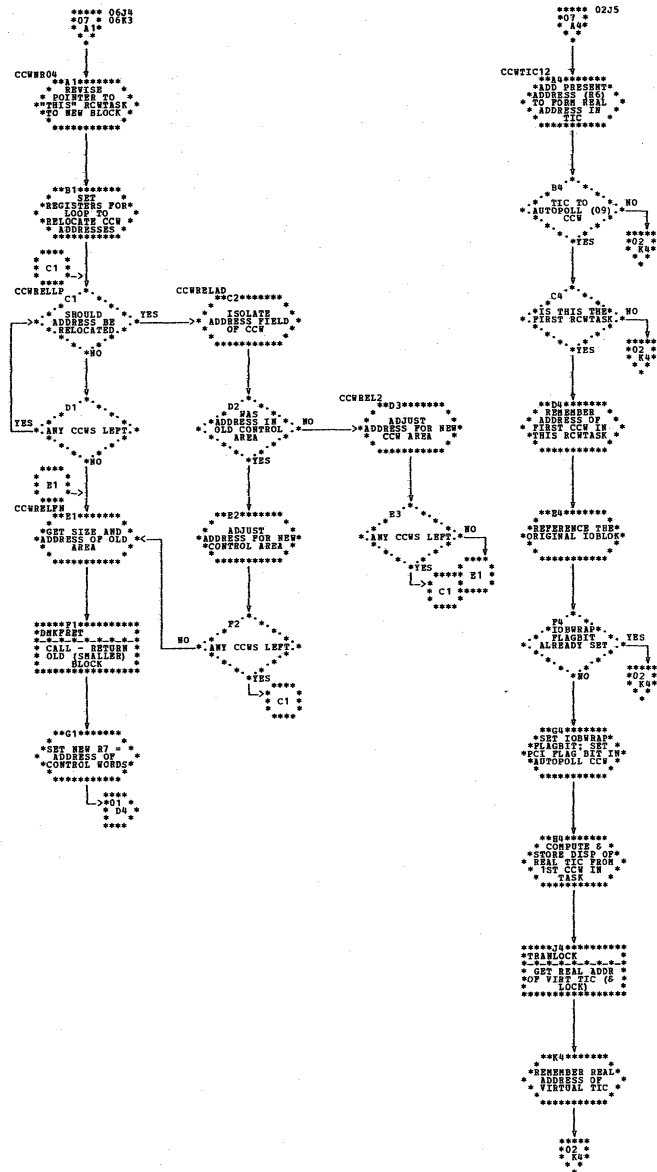




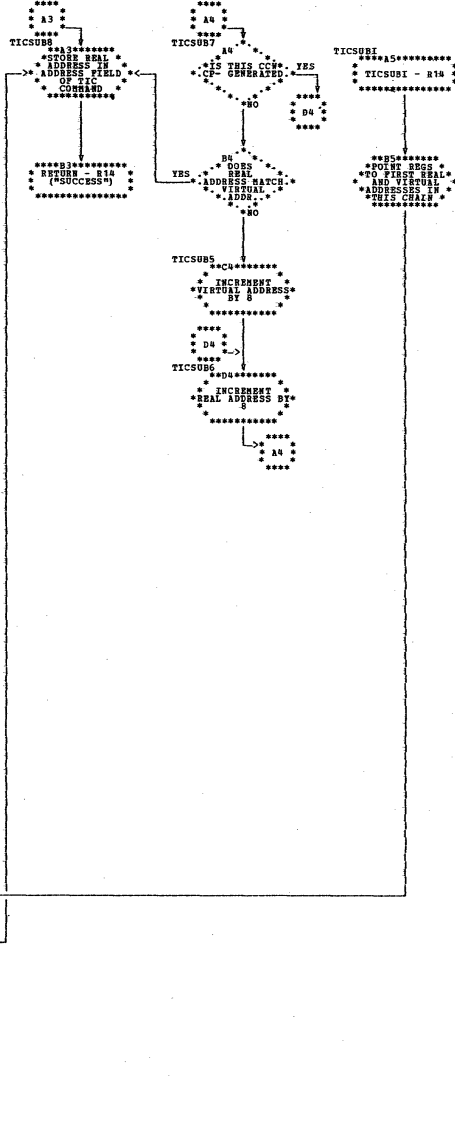
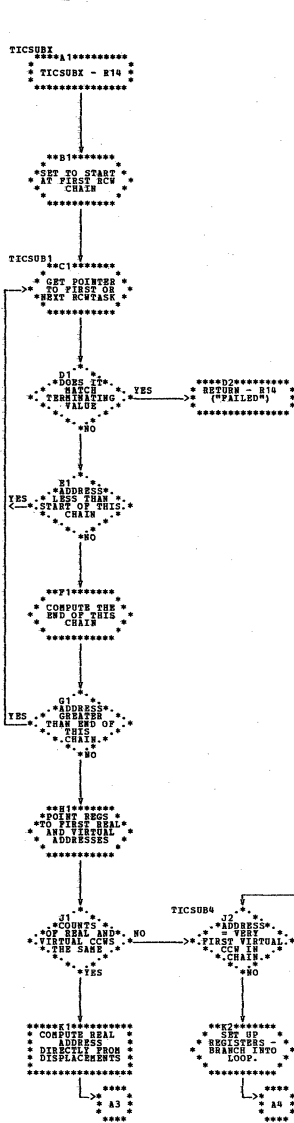


| DMKCCW -- Translate a Virtual CCW List to a Real List (Parts 5 and 6 of 31)

| DMKCCW -- Translate a Virtual CCW List to a Real List (Parts 7 and 8 of 31)

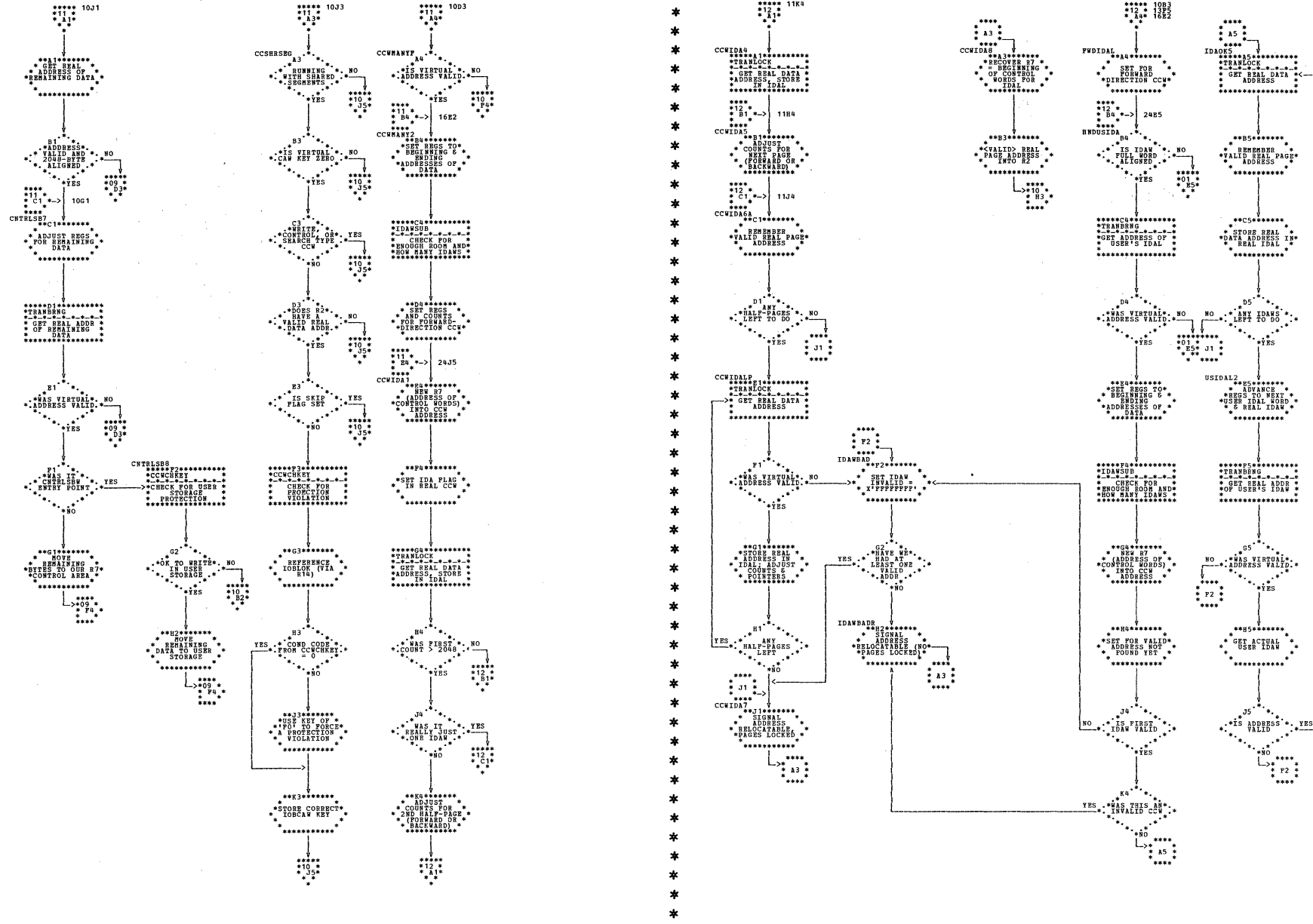


\* \* \* \* \*



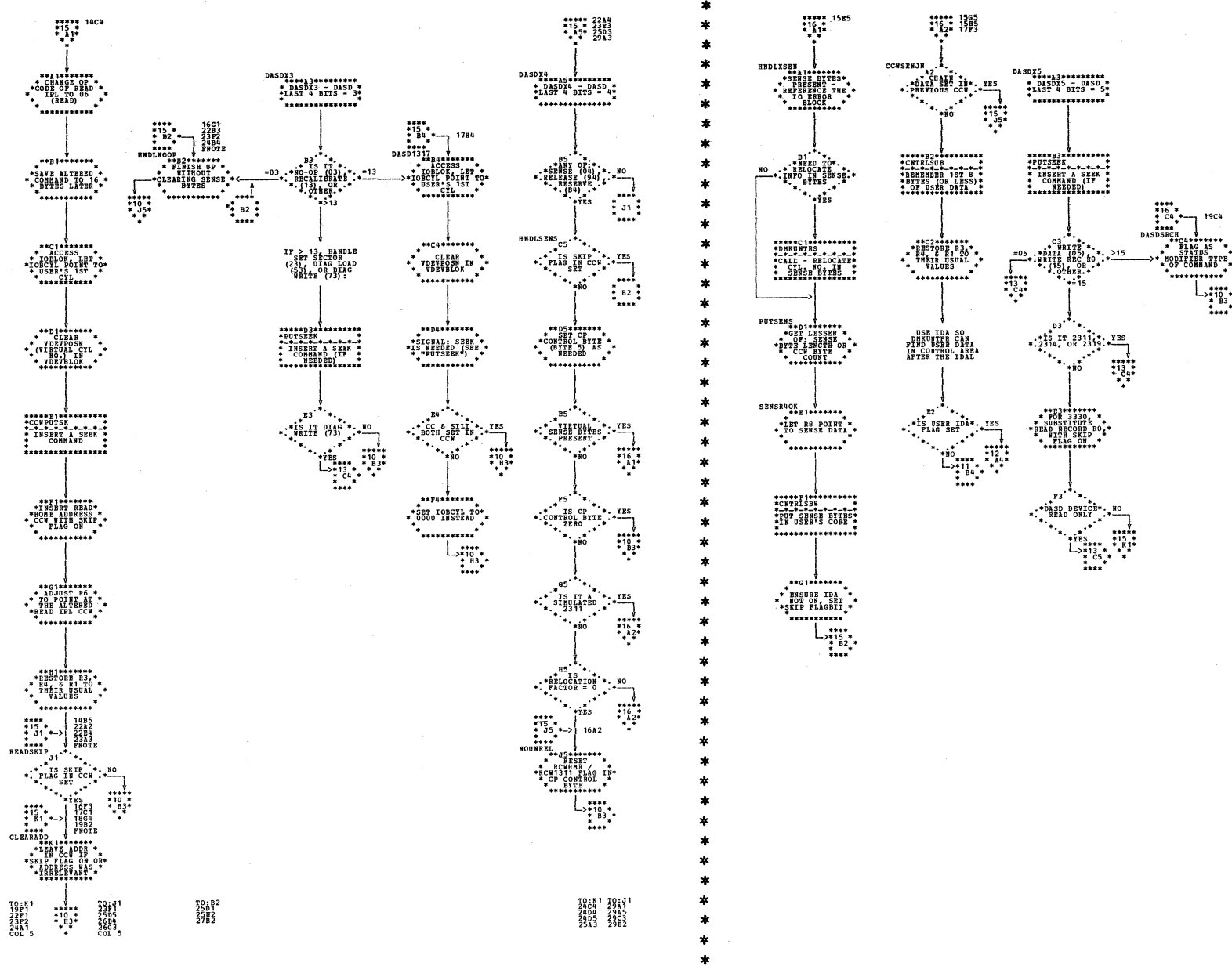


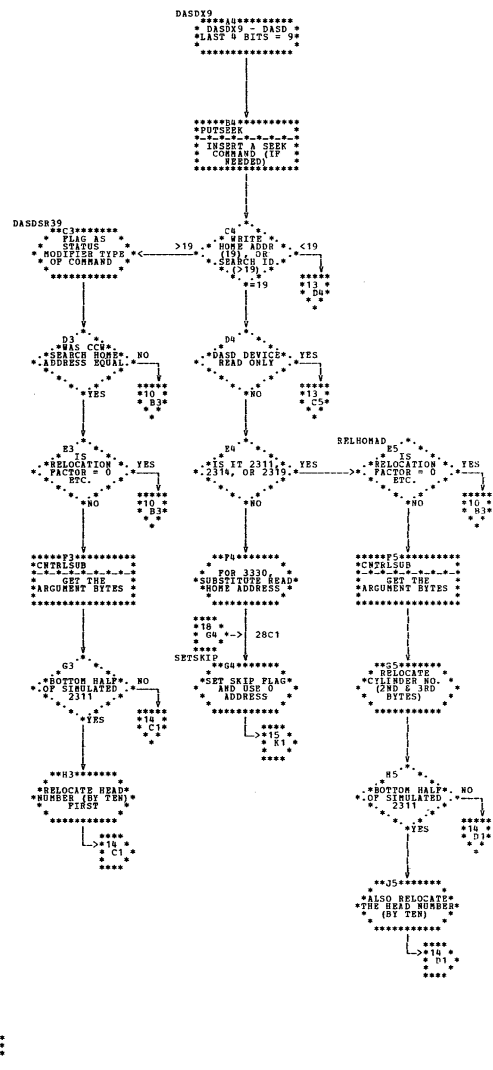
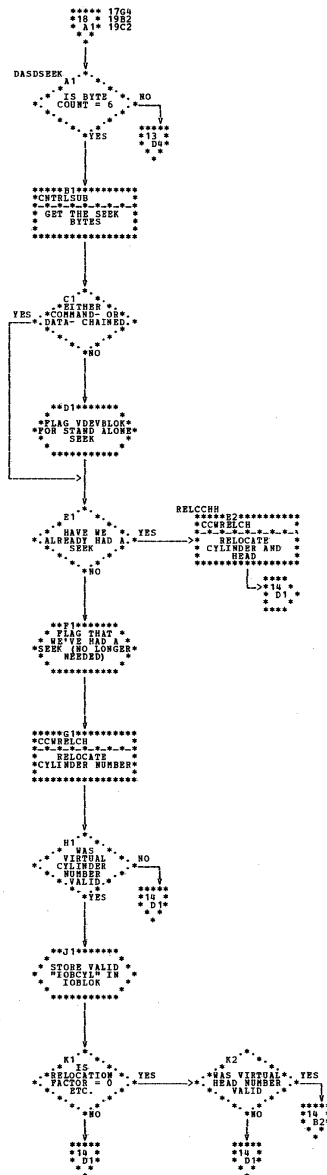
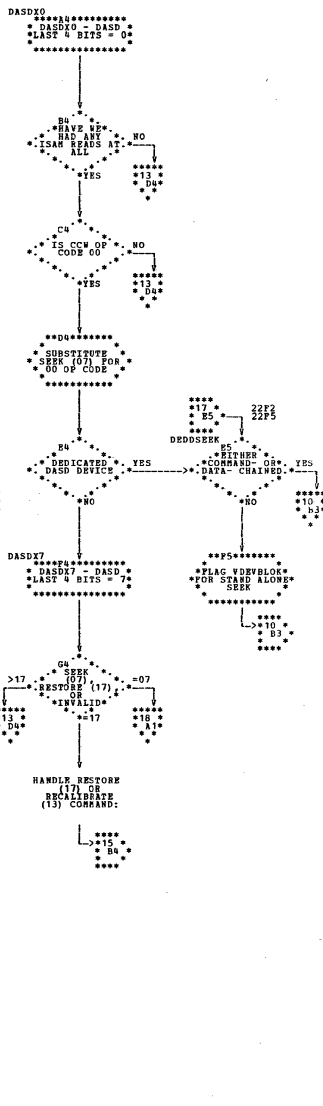
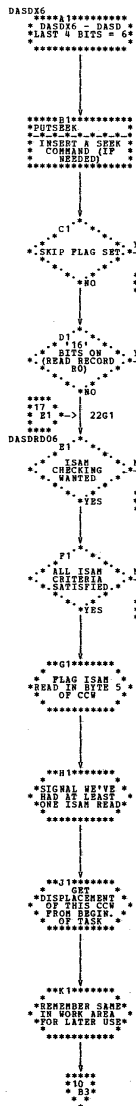
| DMKCCW -- Translate a Virtual CCW List to a Real List (Parts 11 and 12 of 31)





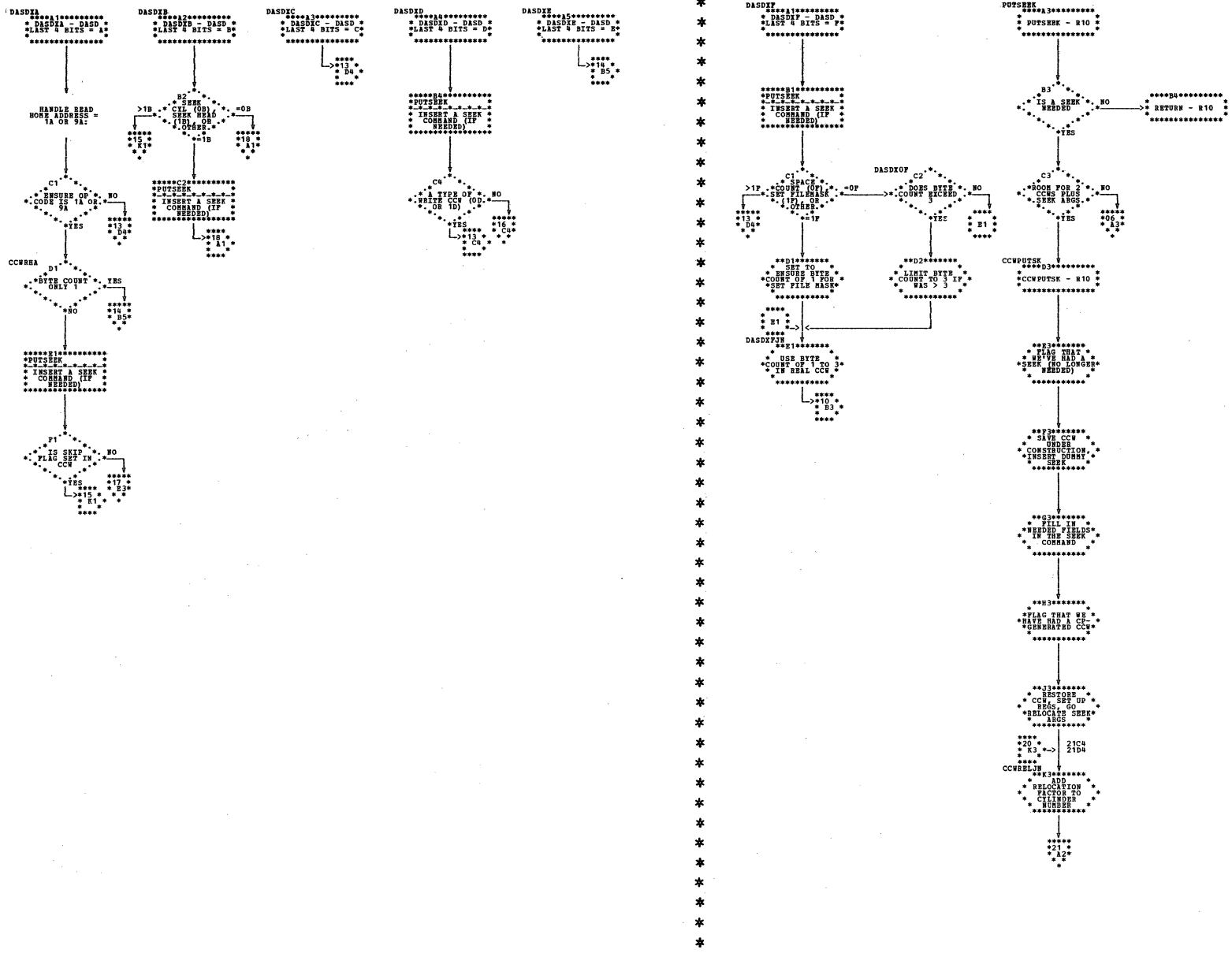
| DMKCCW -- Translate a Virtual CCW List of a Real List (Parts 15 and 16 of 31)





| DMKCCW -- Translate a Virtual CCW List to a Real List (Parts 17 and 18 of 31)

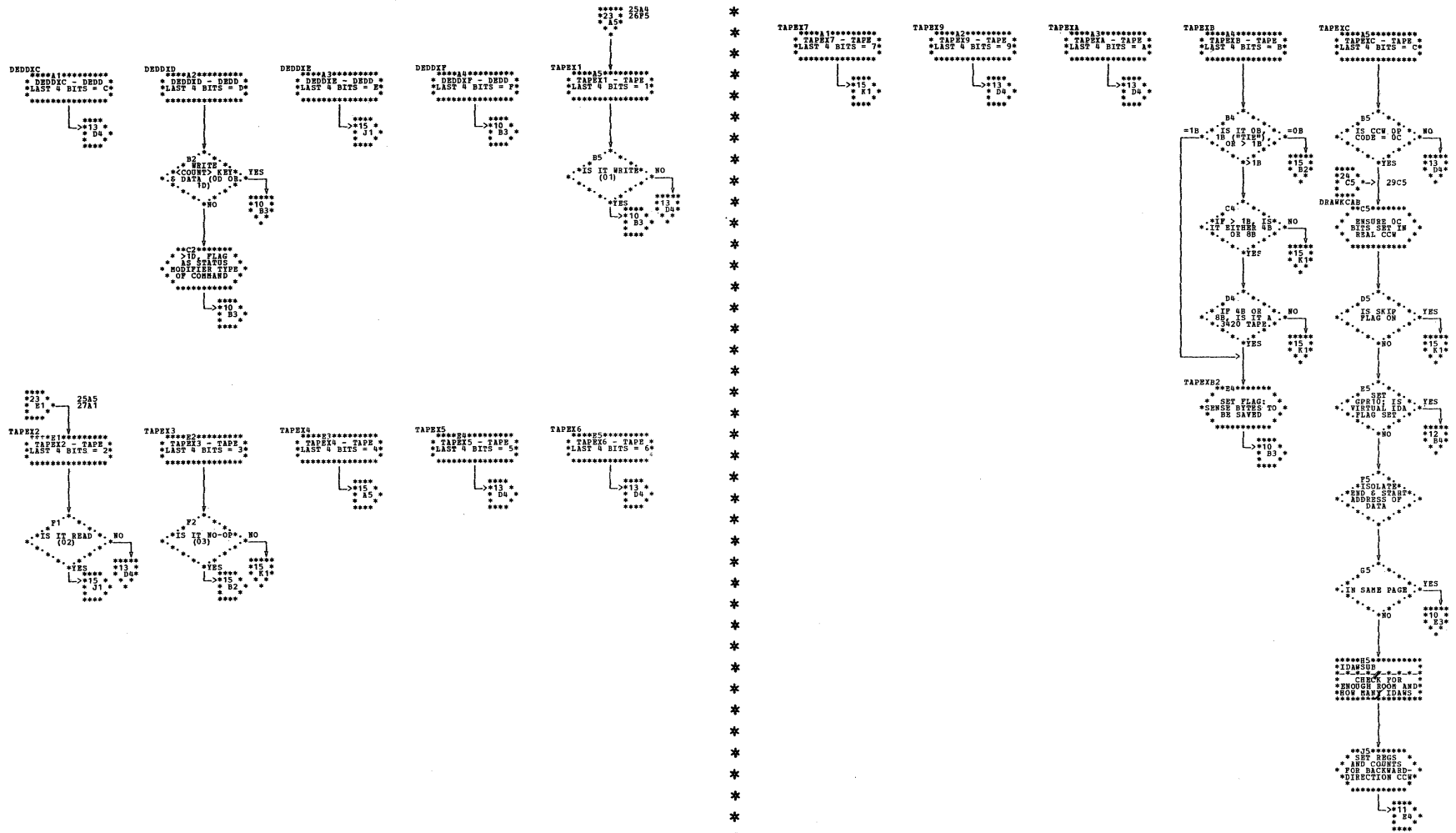
DMKCCW -- Translate a Virtual CCW List to a Real List (Parts 19 and 20 of 31)

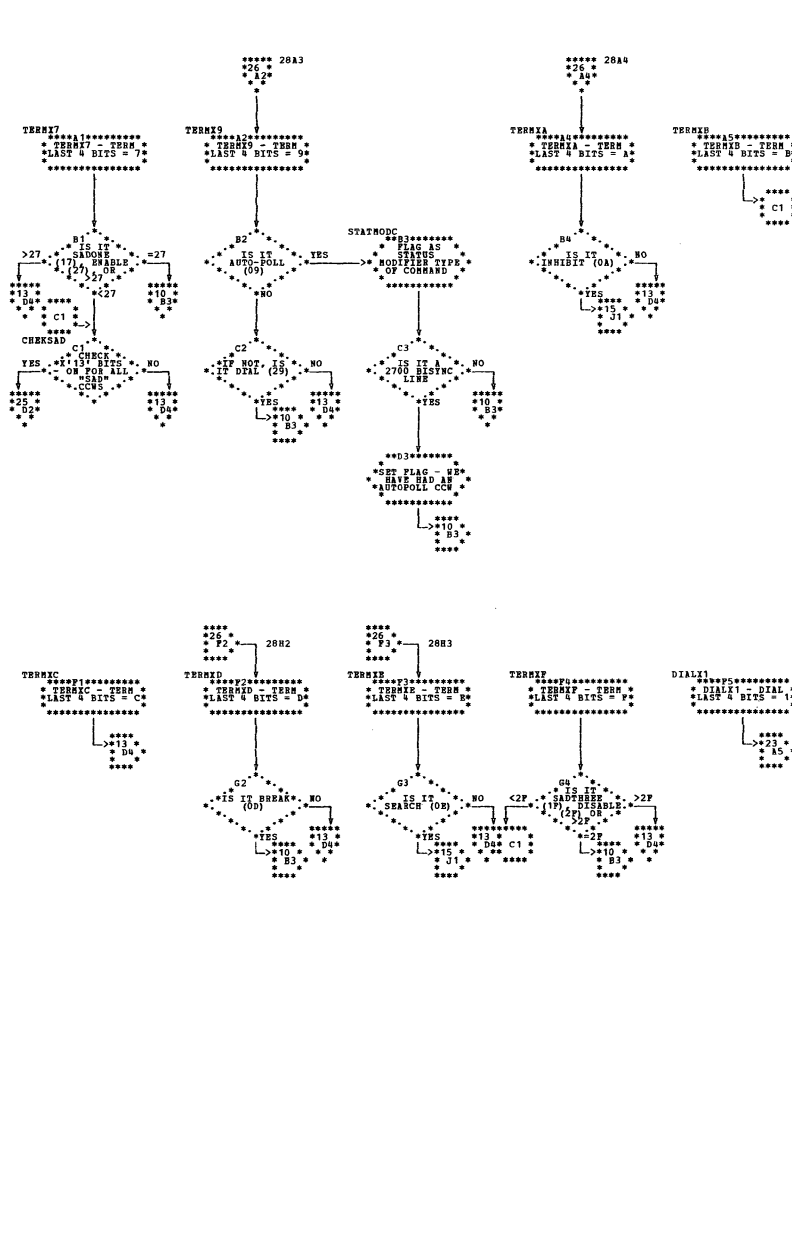
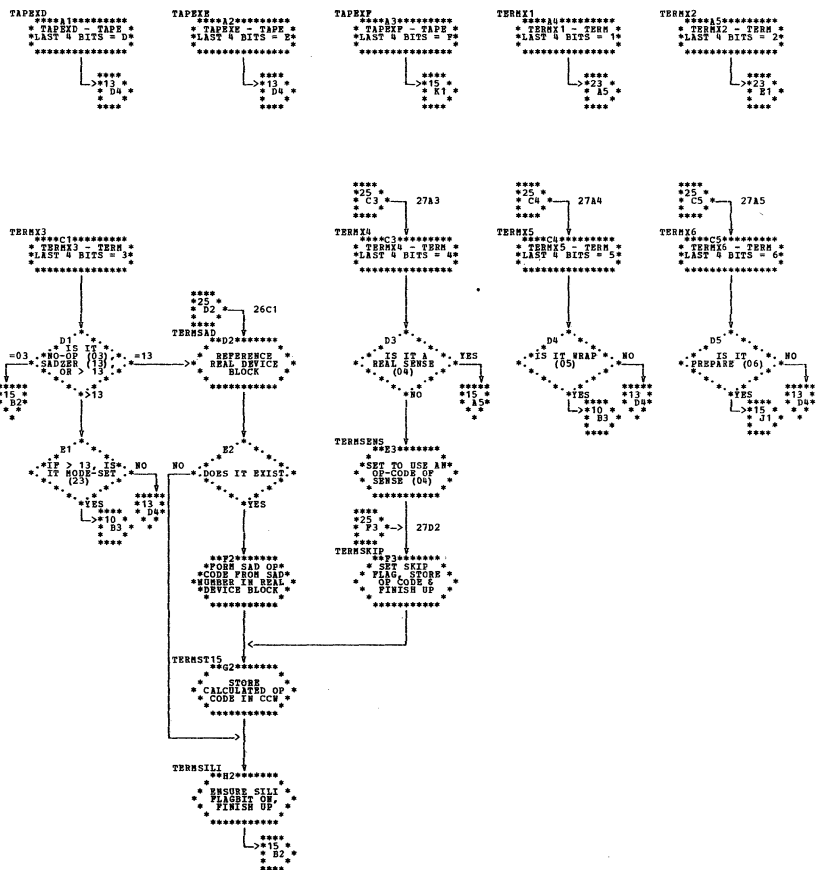






| DMKCCW -- Translate a Virtual CCW List to a Real List (Parts 23 and 24 of 31)



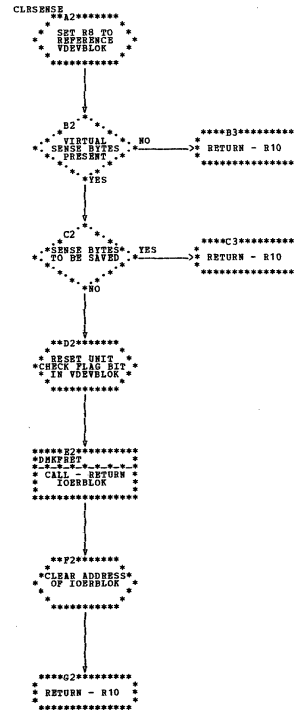


| DMKCCW -- Translate a Virtual CCW List to a Real List (Parts 25 and 26 of 31)





DMKCCW -- Translate a Virtual CCW List to a Real List (Part 31 of 31)



DMKCBDI  
\*\*\*\*\*  
\*DISPLAY VIRTUAL\*  
\*STORAGE\*

\*\*B1\*\*  
\*SET DISPLAY\*  
\*VIRT. FLAG\*

\*\*\*  
\*O1 -> 0385  
\*C1 -> 0481  
\*      0482  
\*\*\*\*\*  
DISCDBI  
\*\*W2\*\*  
\*\*\*\*\*  
\*CALL DMPRES  
\*GET DUMBUF  
\*BUFFER\*

\*\*D1\*\*  
\*SET UP LINK\*  
\*LENGTH AND\*  
\*TRANSLATE\*  
\*COUNT\*

\*\*\*  
\*O1 -> 03A3  
\*D1 -> 11E2  
\*      11E3  
\*\*\*\*\*  
DISCDBI  
\*\*W1\*\*  
\*\*\*\*\*  
\*DISWRITE  
\*WRITE BUFFER  
\*AND REINITIALIZE\*

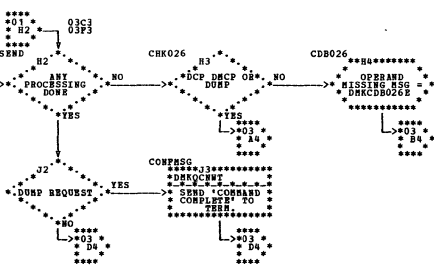
\*ALREADY  
\*HAVE ARGUMENT\*

DISCDBI  
\*\*W1\*\*  
\*\*\*\*\*  
\*GET NEXT  
\*ARGUMENT\*

\*ARGUMENT  
\*FOUND\*

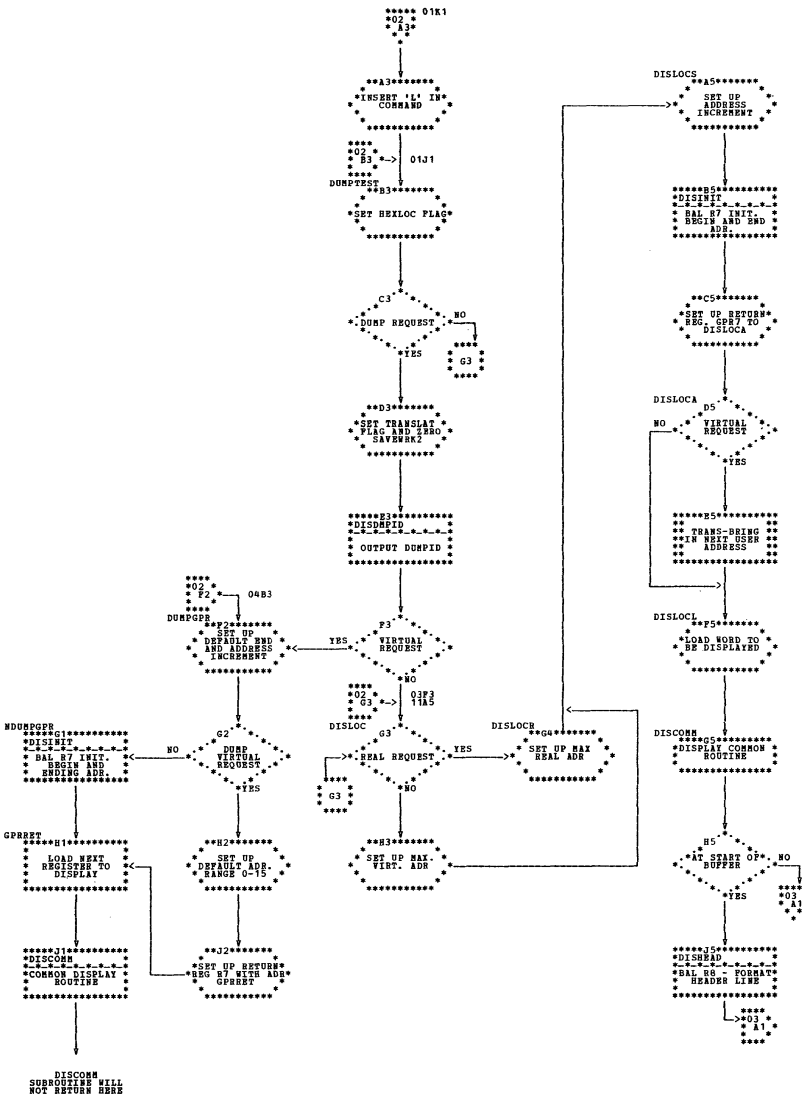
\*DISPLAY TYPE\*

PCB	-> 11A0
S	-> 09A3
L	-> 02B3
E	-> 11E3
L	-> 05A2
CAL	-> 11G3
I	-> 00E3

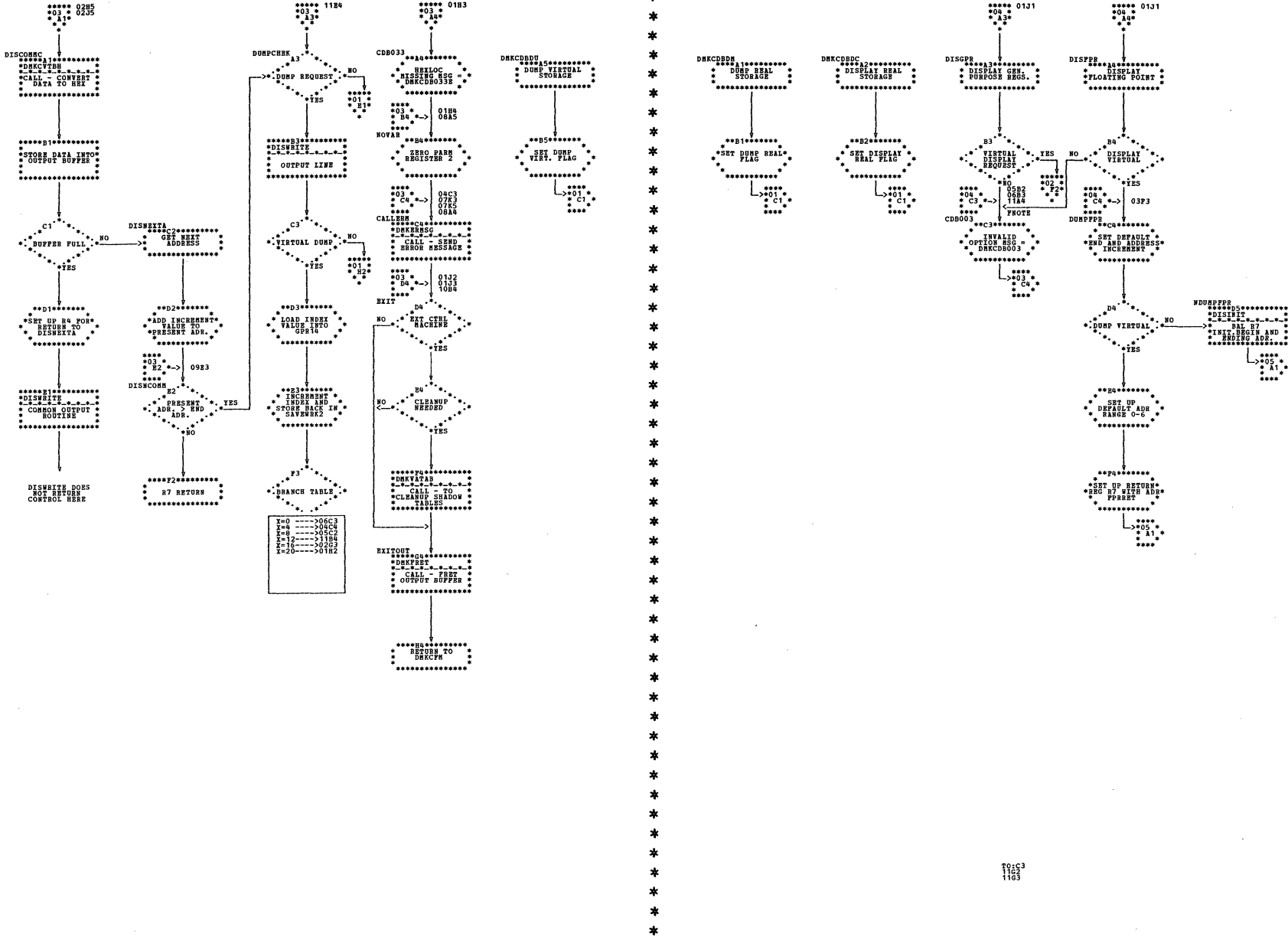


DMKCDI -- Process DCP, DISPLAY, DMCP and DUMP commands (Parts 1 and 2 of 11)

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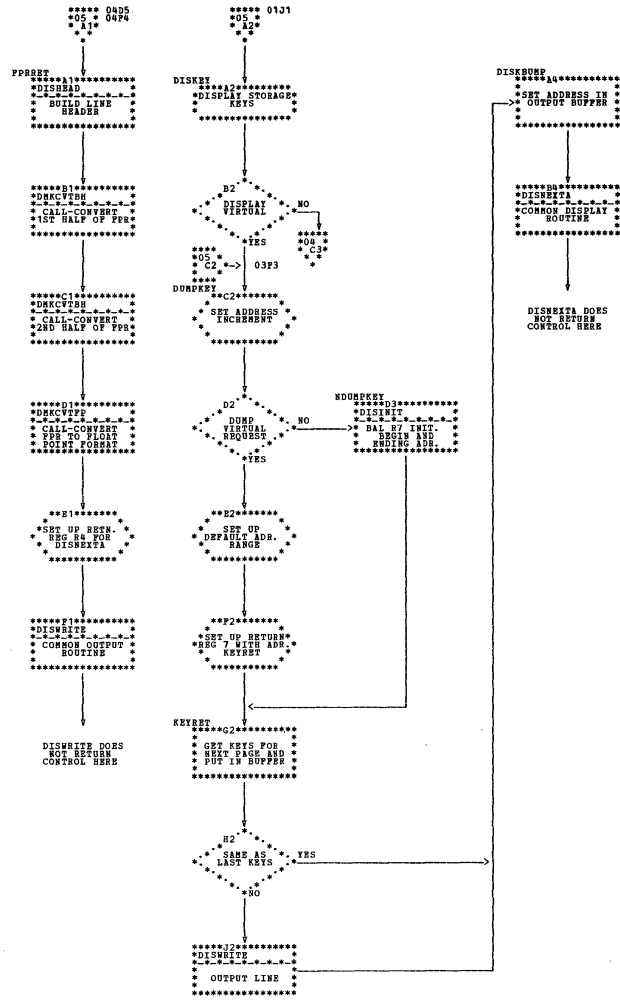


DMKDCB -- Process DCP, DISPLAY, DMCP and DUMP Commands (Parts 3 and 4 of 11)

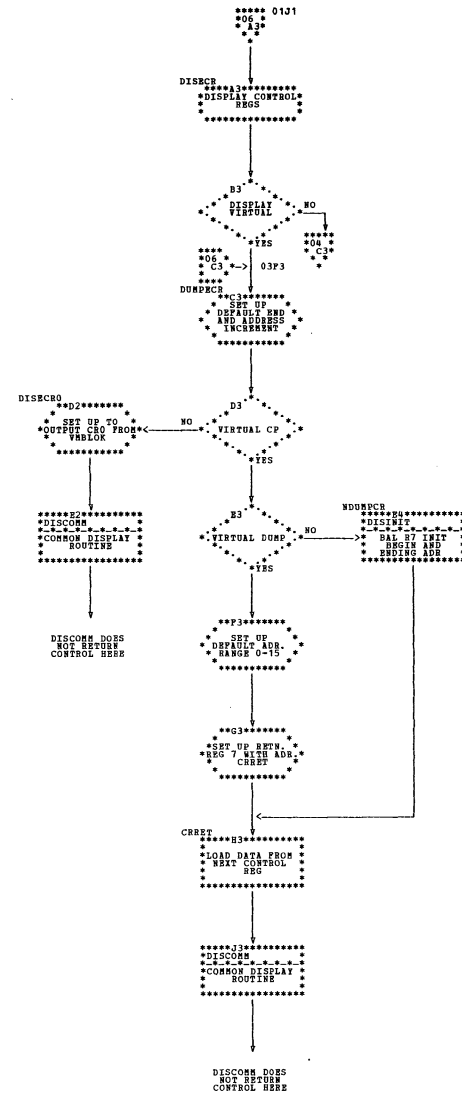


01C3  
1123  
1163

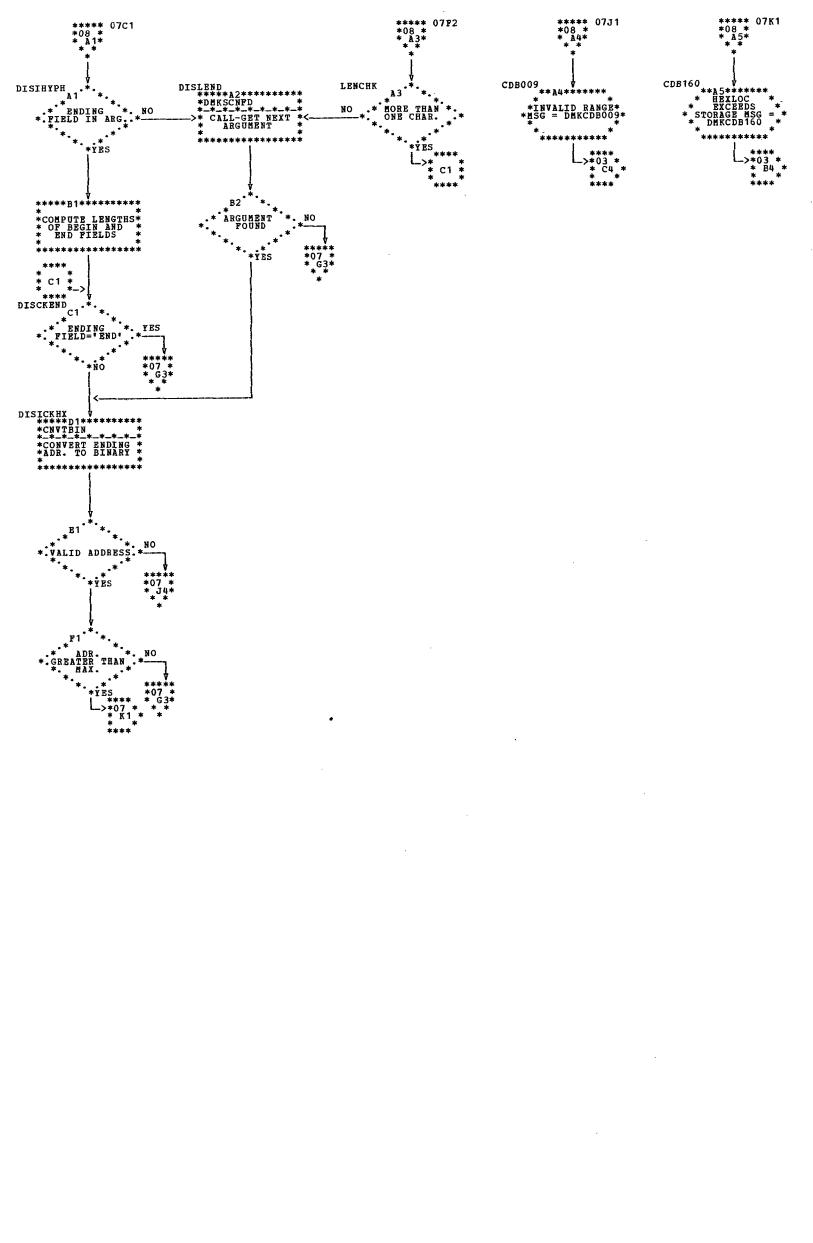
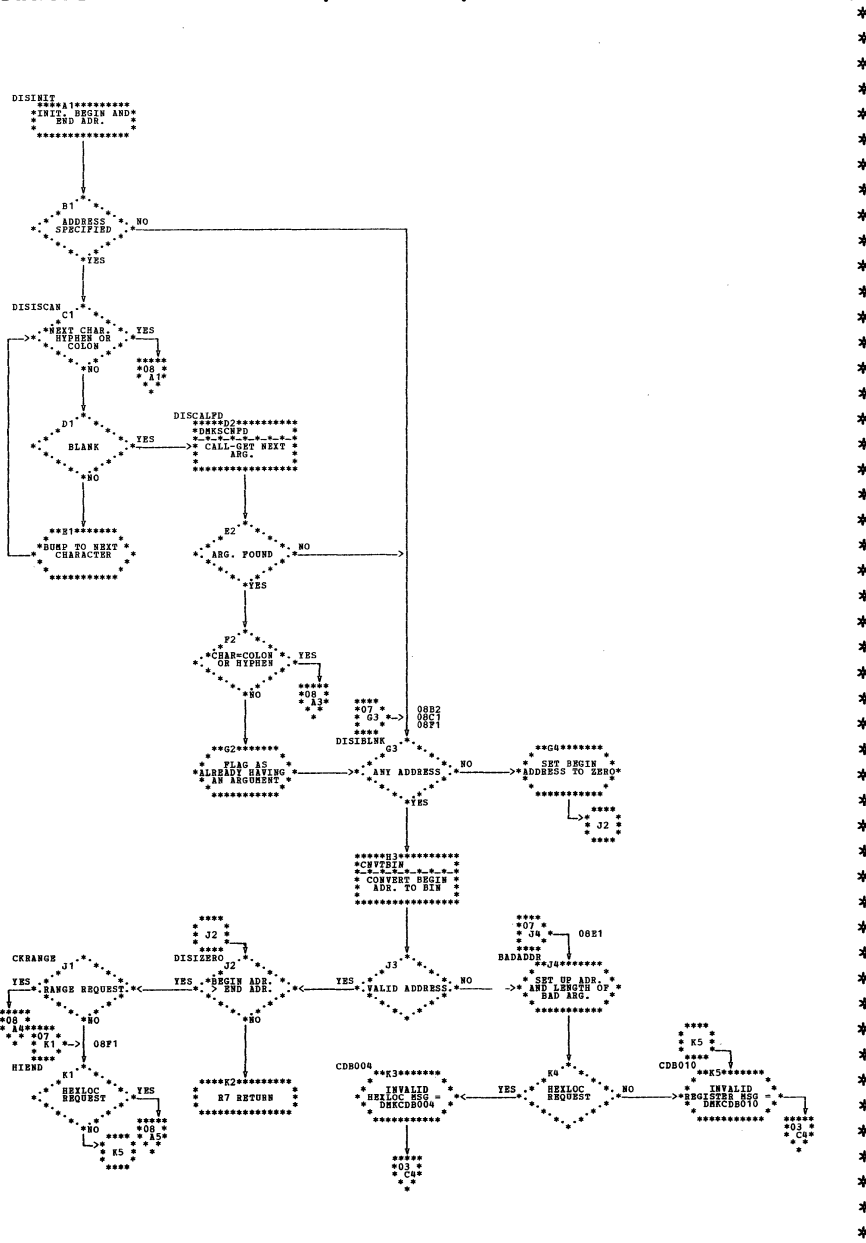




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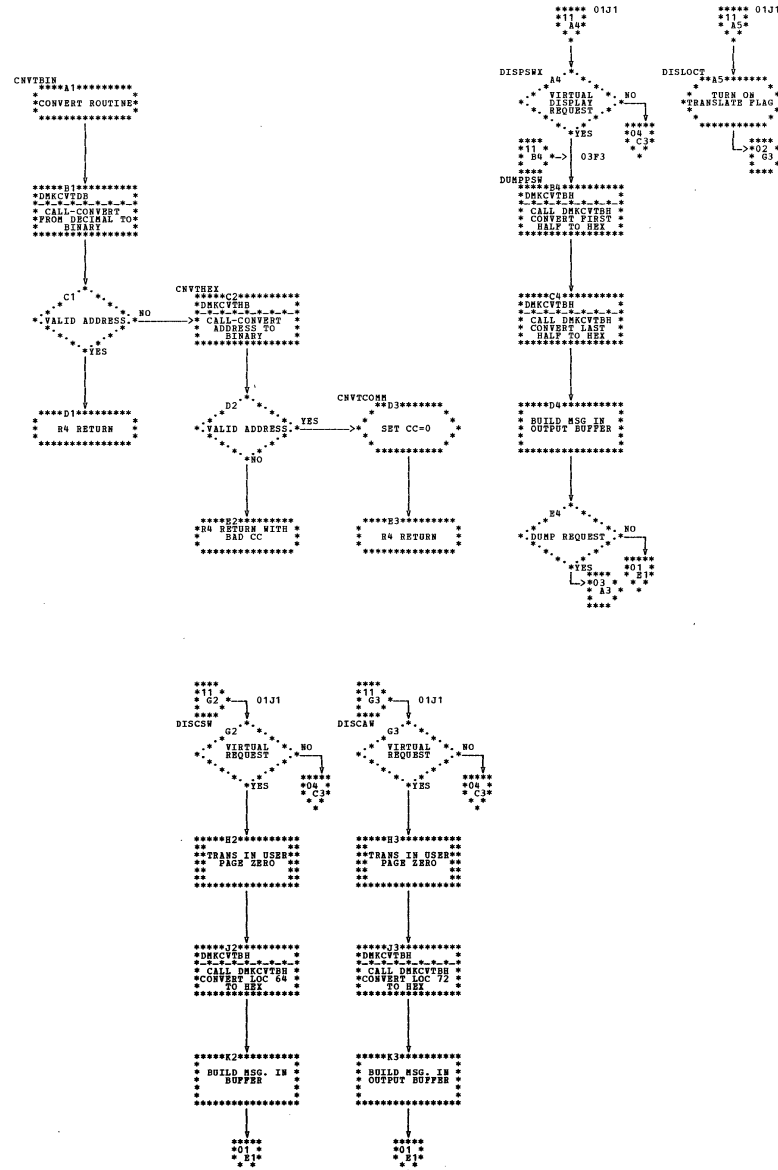


DMKCIDB -- Process DCP, DISPLAY, DMCP and DUMP Commands (Parts 7 and 8 for 11)



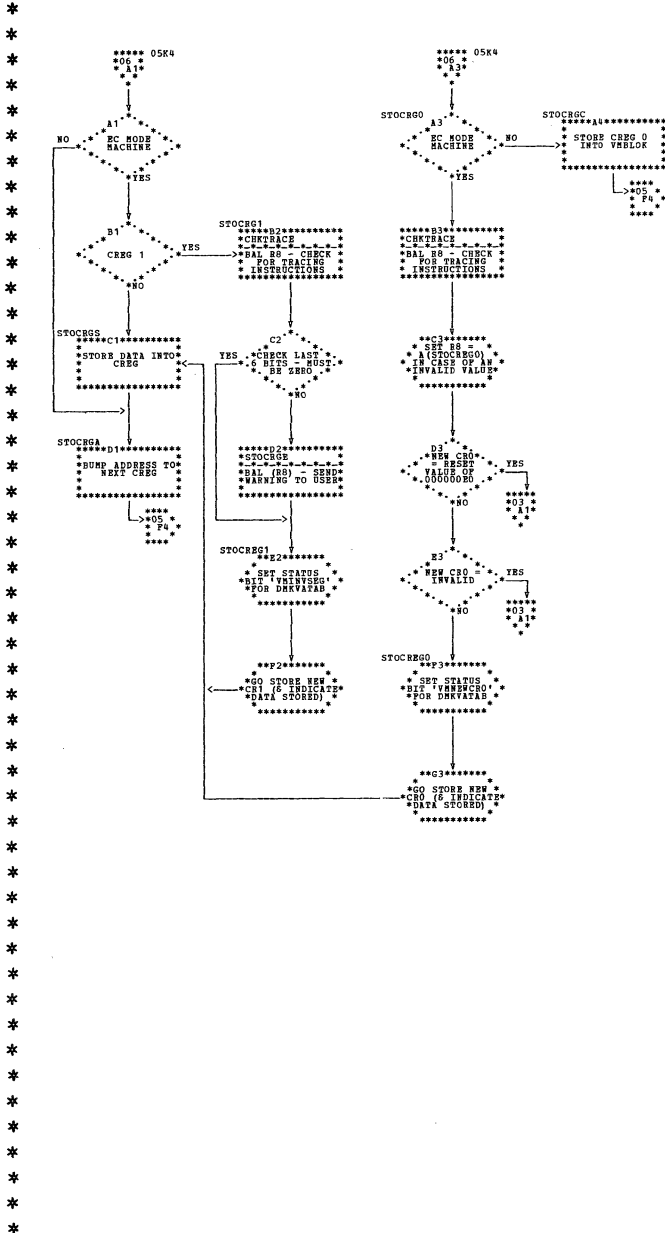
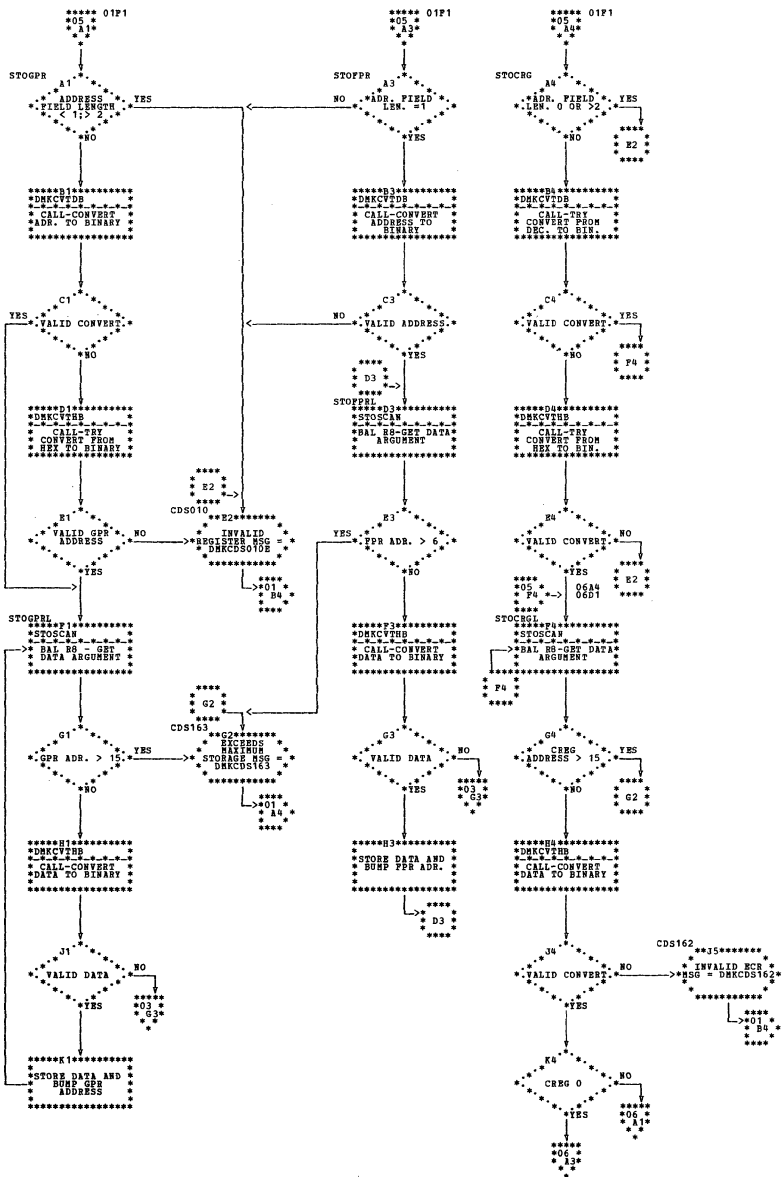


DMKCDB -- Process DCP, DISPLAY, DMCP and DUMP Commands (Part 11 of 11)



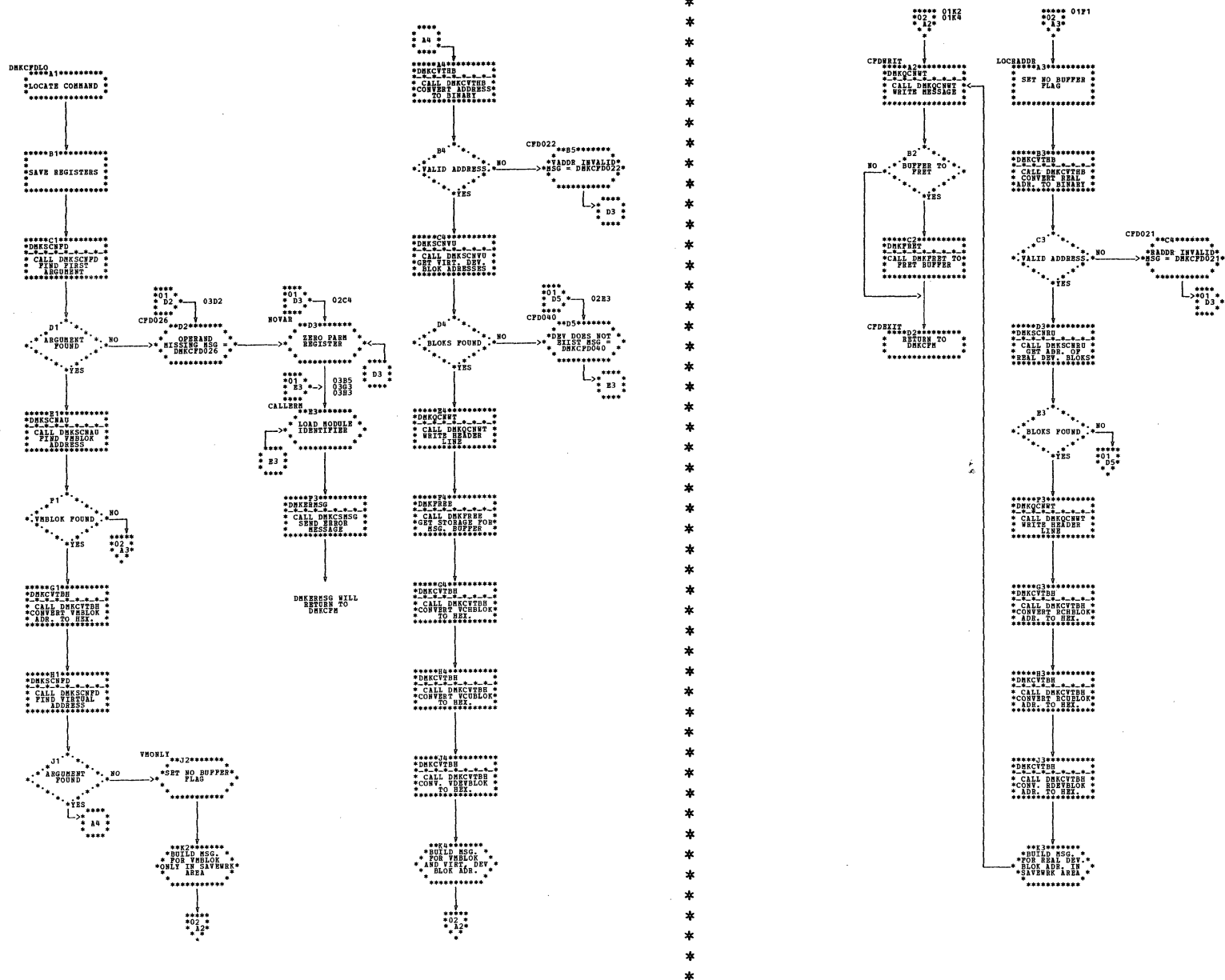






| DMKCD5 -- Process STCEP and STORE Commands (Parts 5 and 6 of 6)

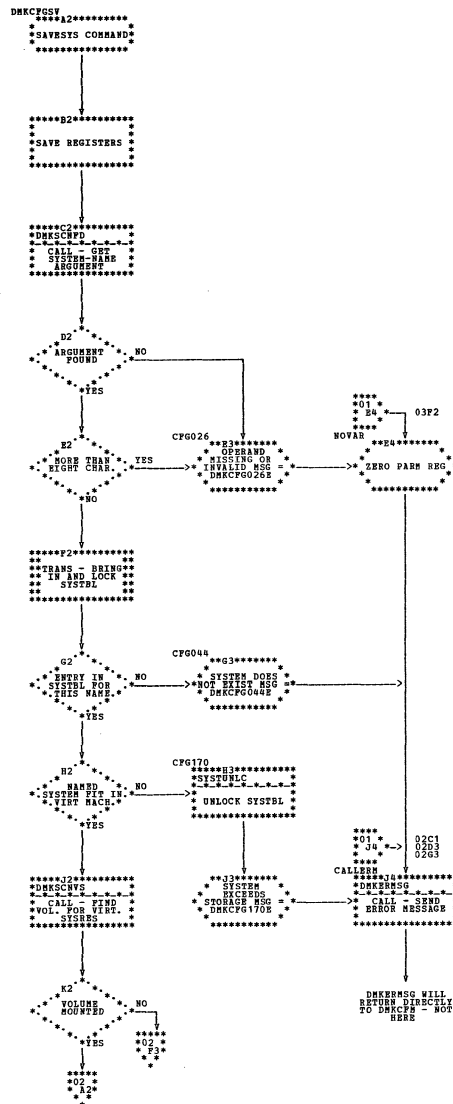
DMKCFD -- Process ADSTOP and LOCATE Commands (Parts 1 and 2 of 4)



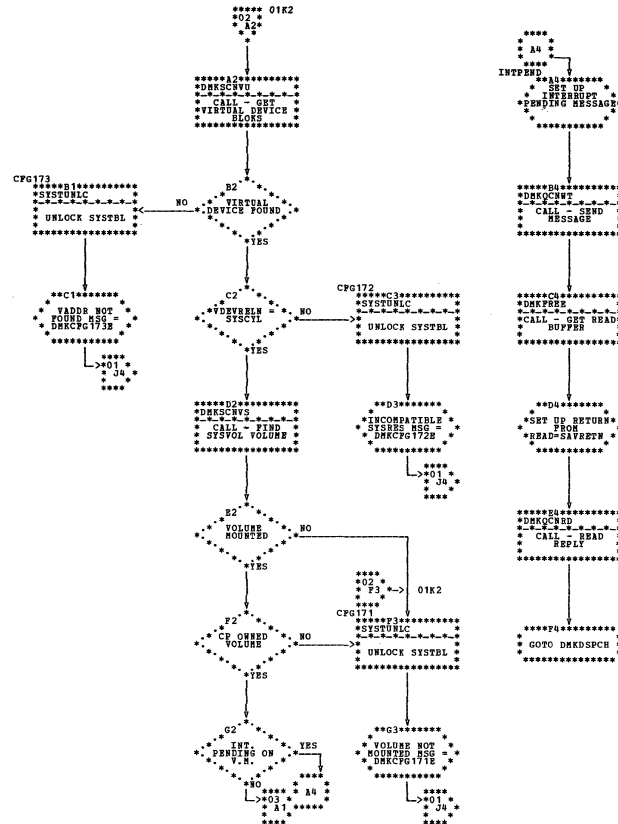


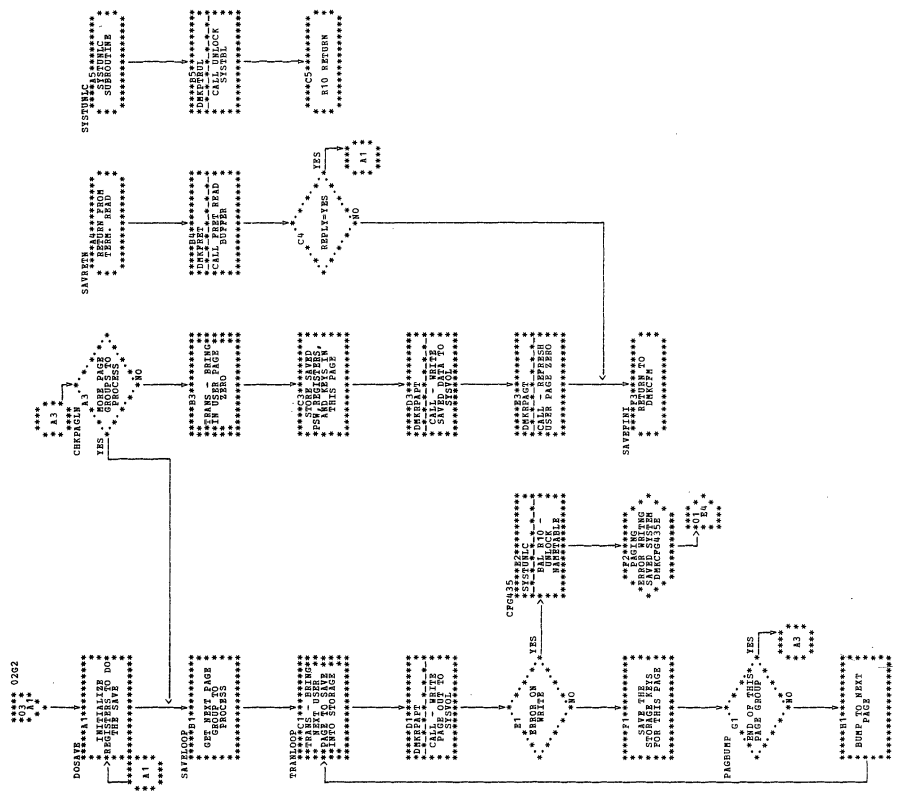


DMKCFG -- Process SAVESYS Command (Parts 1 and 2 of 3)

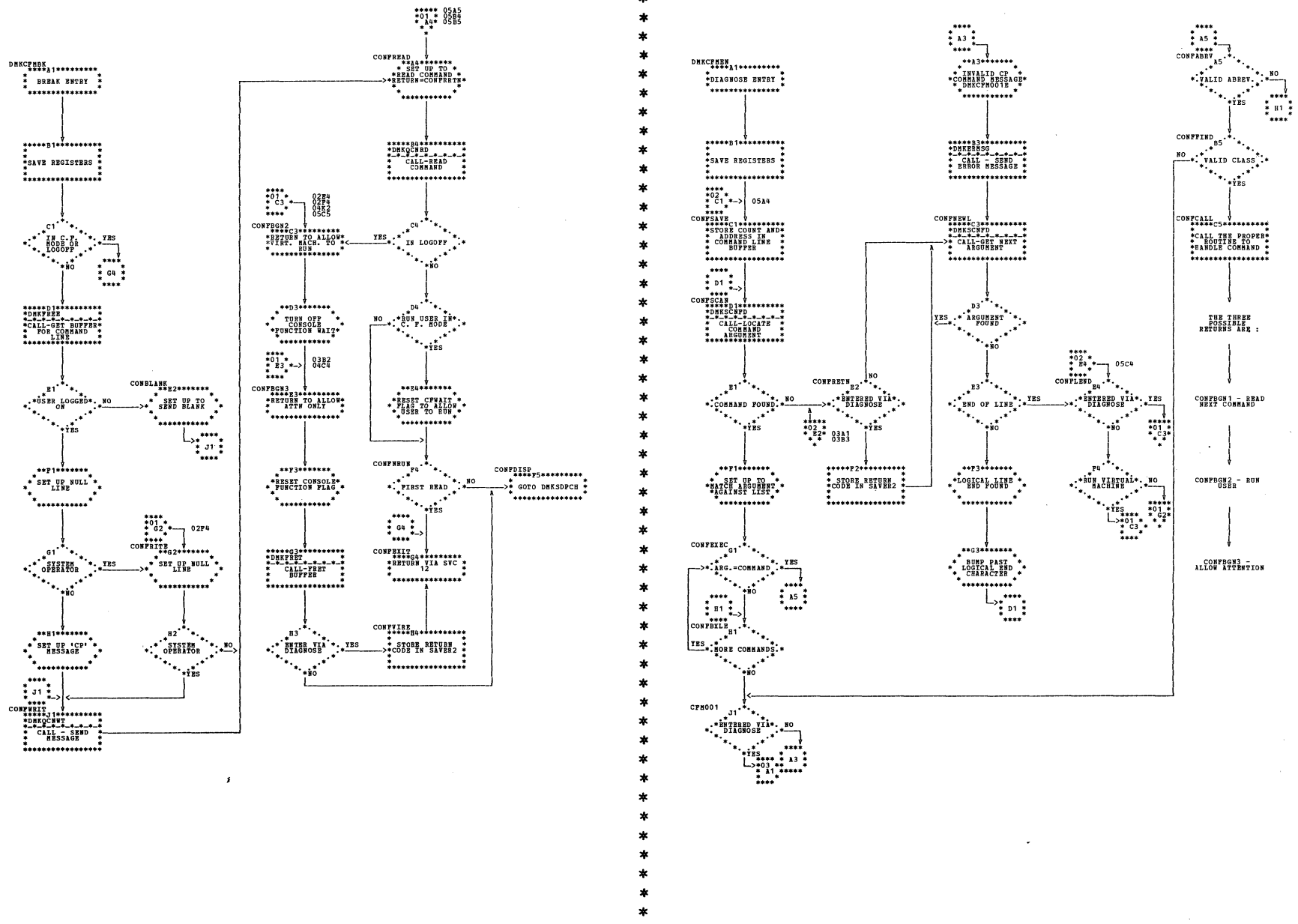


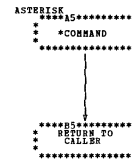
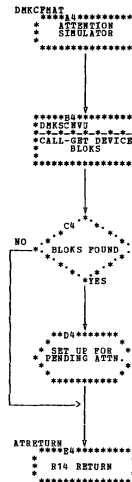
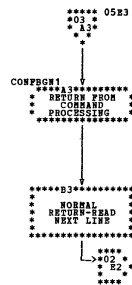
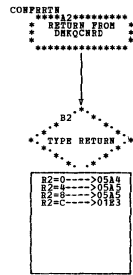
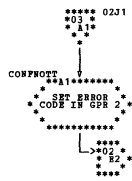
\* \* \* \* \*



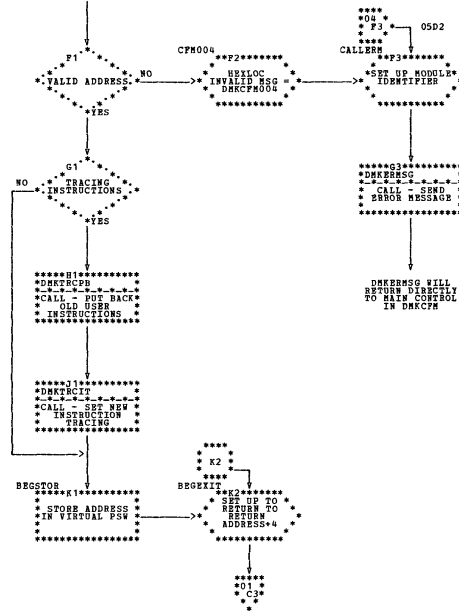
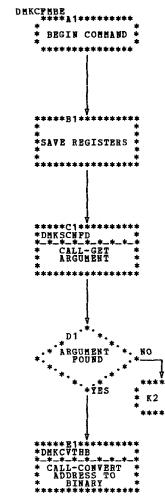


DMKCFM -- Main Console Function Routine (Parts 1 and 2 of 5)



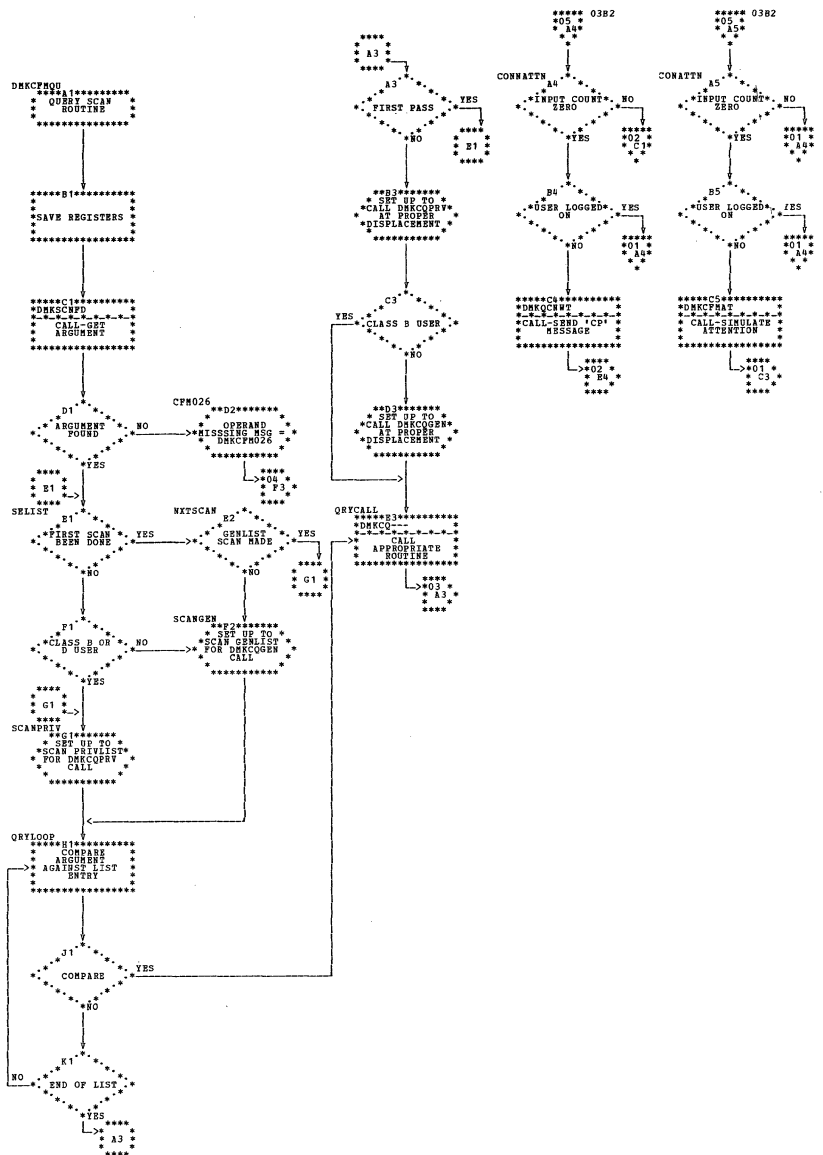


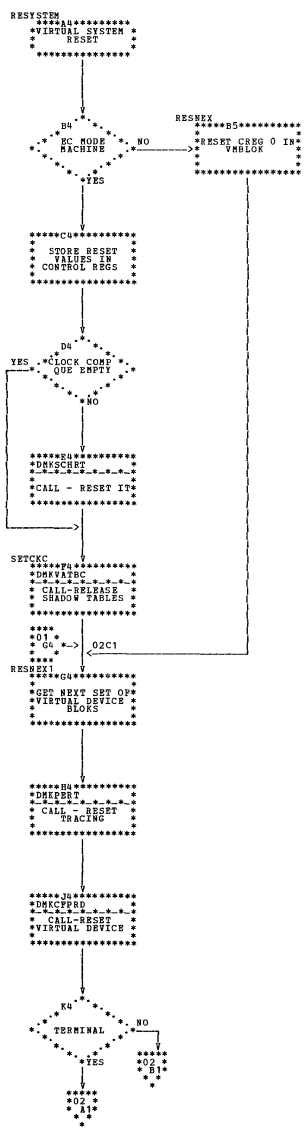
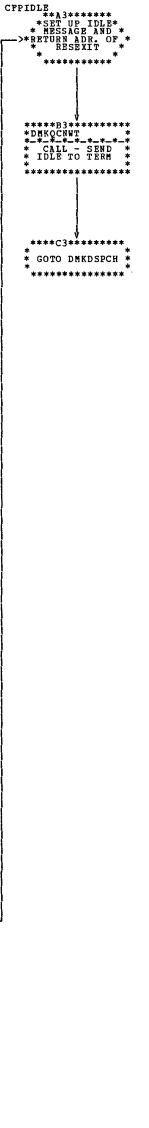
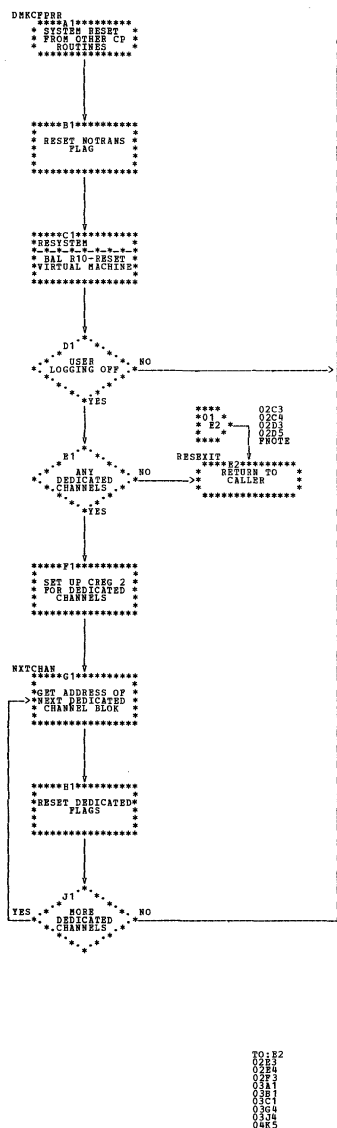
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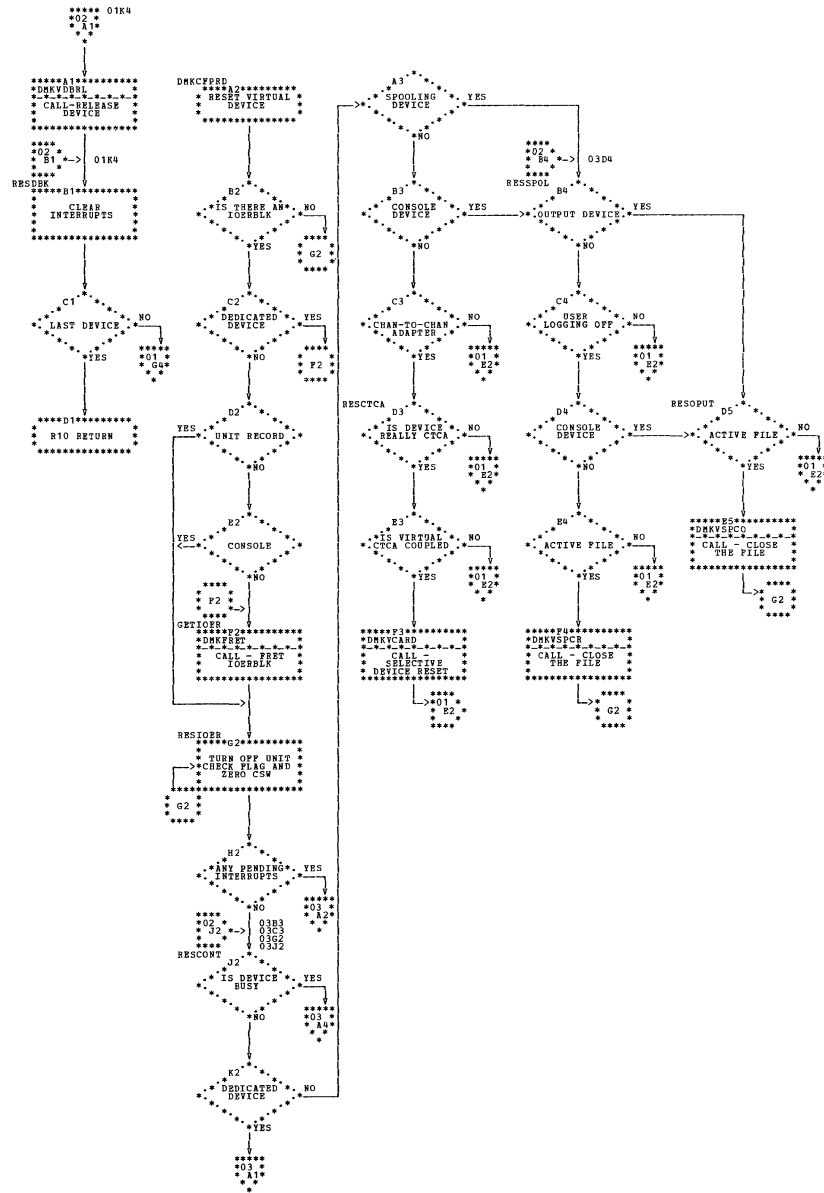
| DMKCFM -- Main Console Function Routine (Parts 3 and 4 of 5)

| DMKCFM -- Main Console Function Routine (Part 5 of 5)





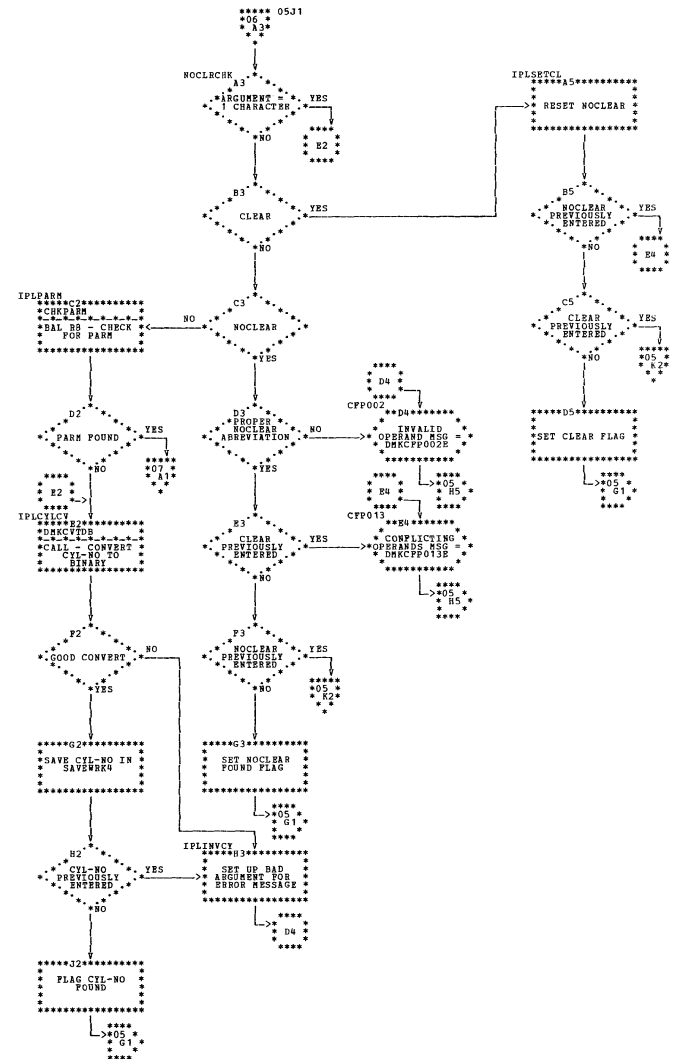
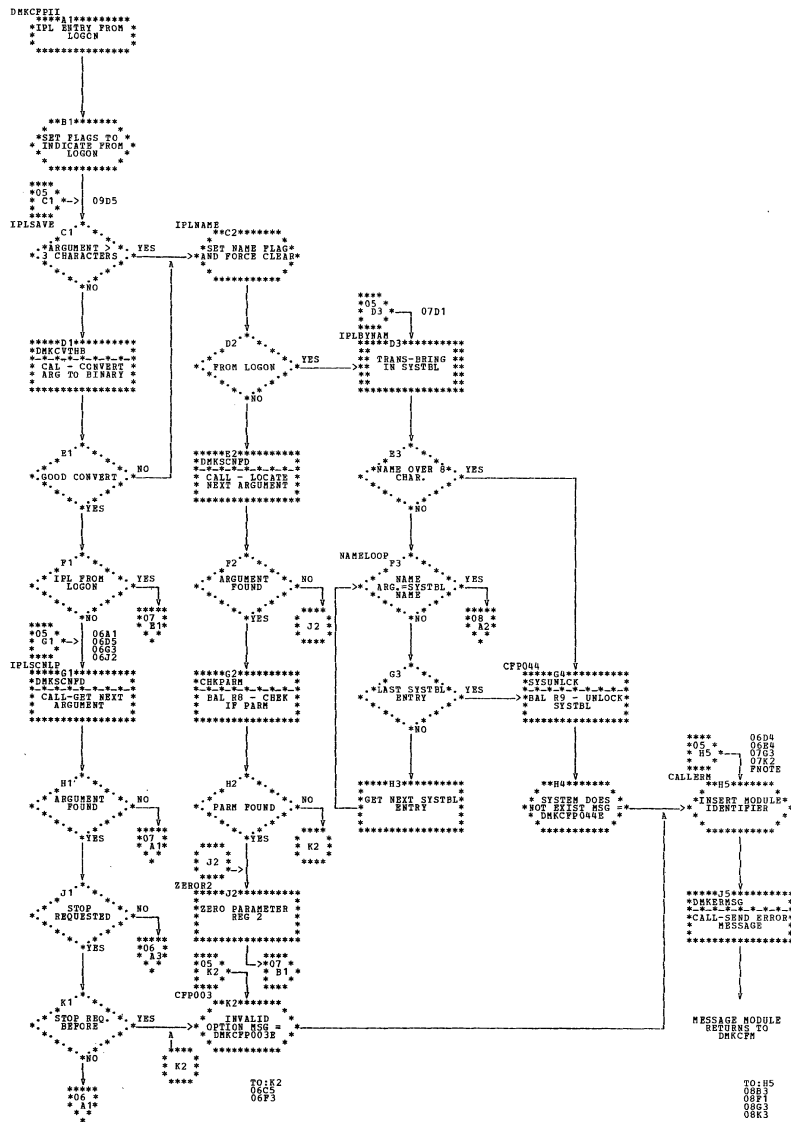
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[ DMKCFP -- Simulate CPU Console to Virtual Machine (Parts 1 and 2 of 10)

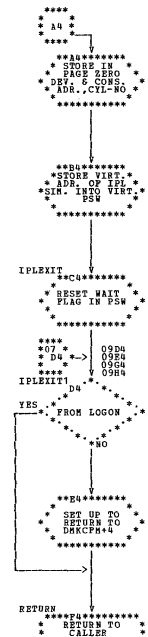
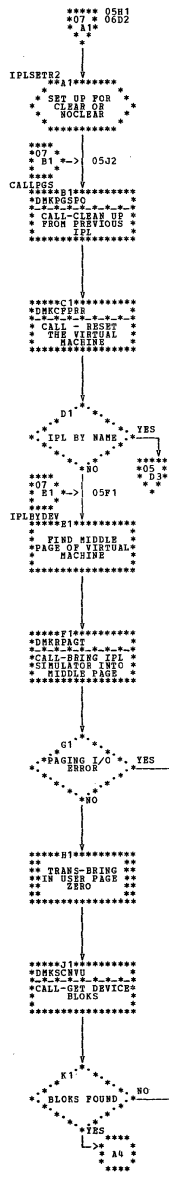




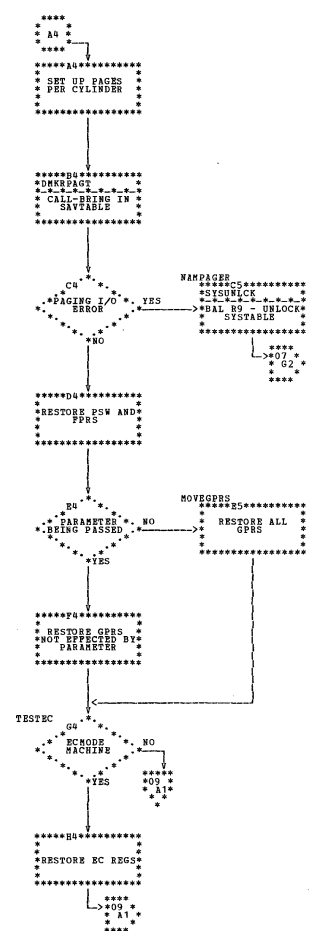
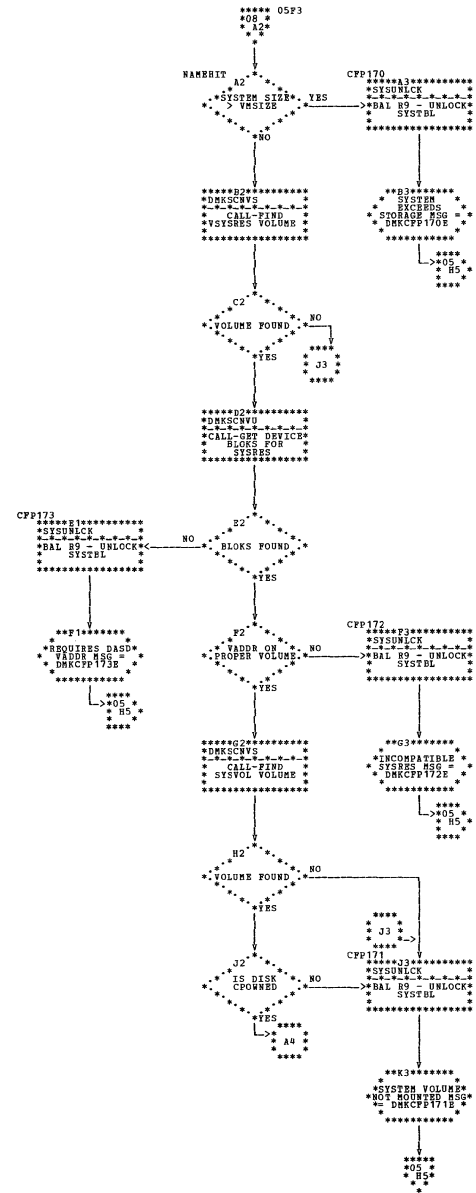


| DMKCFP -- Simulate CPU Console to Virtual Machine (Parts 5 and 6 of 10)

| DMKCFP -- Simulate CPU Console to Virtual Machine (Parts 7 and 8 of 10)

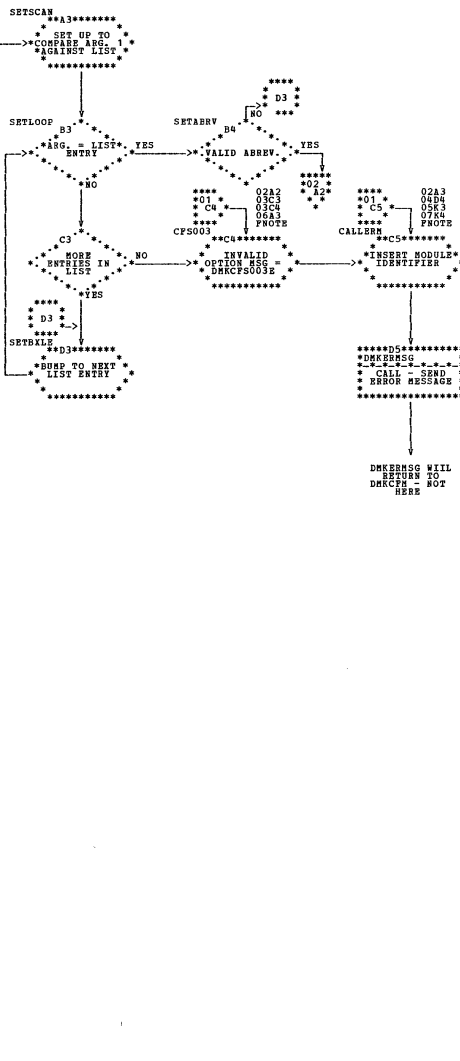
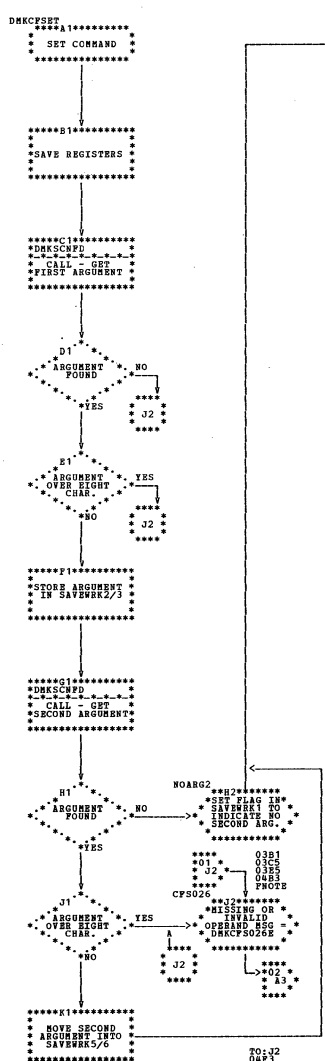


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| DMKCF5 -- Process SET Command (Parts 1 and 2 of 14)

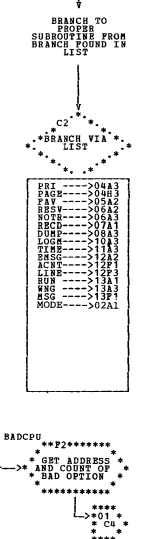
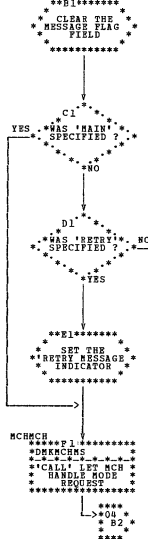
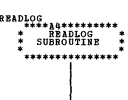
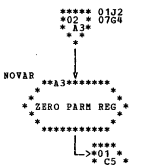
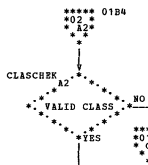


TO: J2  
04F1  
04B5  
04C3  
05A7  
05B2  
06C1  
08A3  
08E7  
COL 5

TO: C4  
09C1  
12C4  
14B3  
14B5

TO: C5 TO: J2  
09D1 10D3  
09D4 11C4  
10E7  
13F5

\* \* \* \* \*

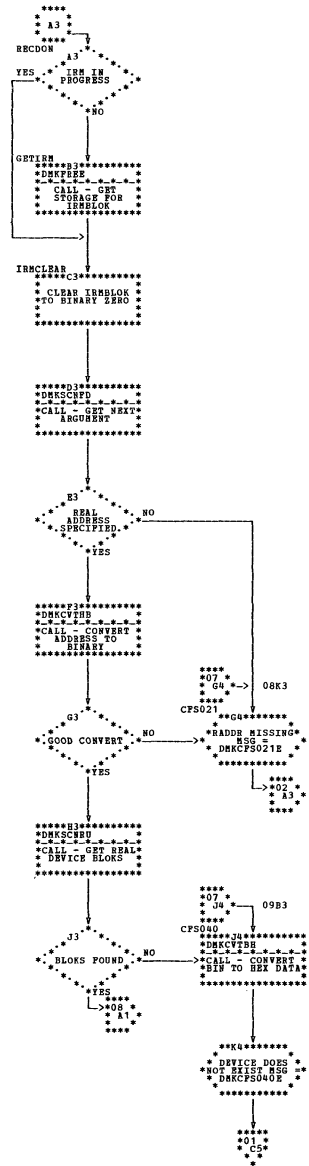


SET	>>08A3
PAGE	>>08B3
REP	>>08A2
RECV	>>06A2
NOB	>>05A3
RECD	>>07A1
UR	>>08A3
LOGH	>>10A3
LINE	>>11A1
RMSG	>>12A2
ACH	>>13A1
LINE	>>14F3
SUB	>>15A1
WNG	>>13A3
SS	>>13F1
MODE	>>02A1

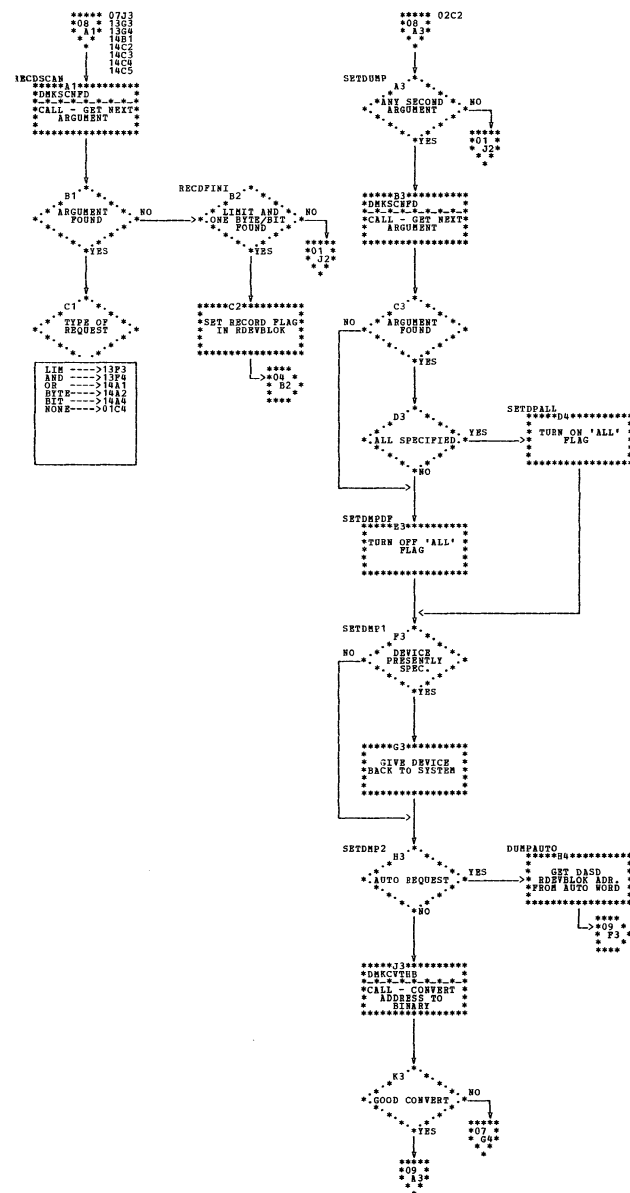








\* \* \* \* \*

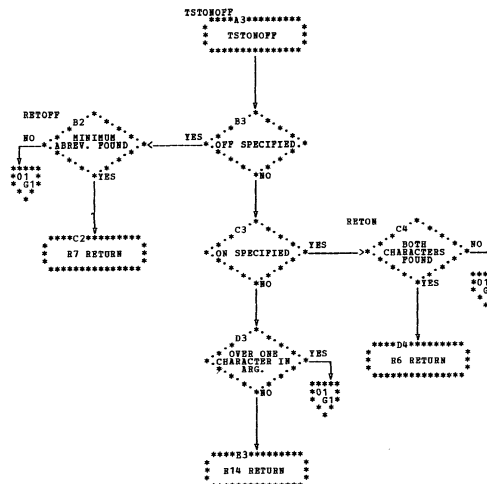
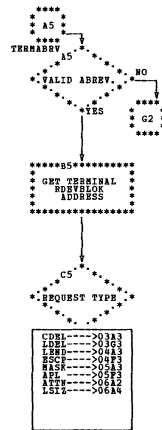
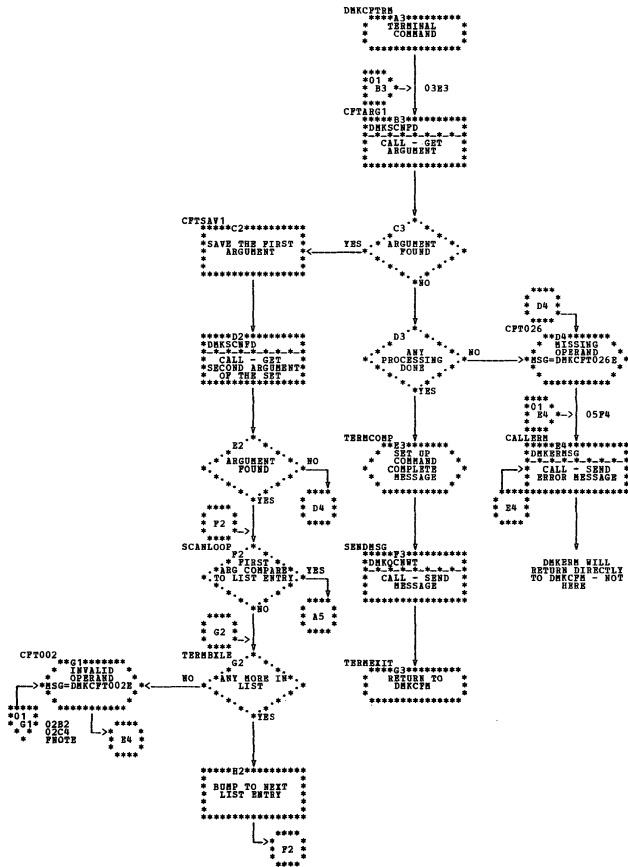






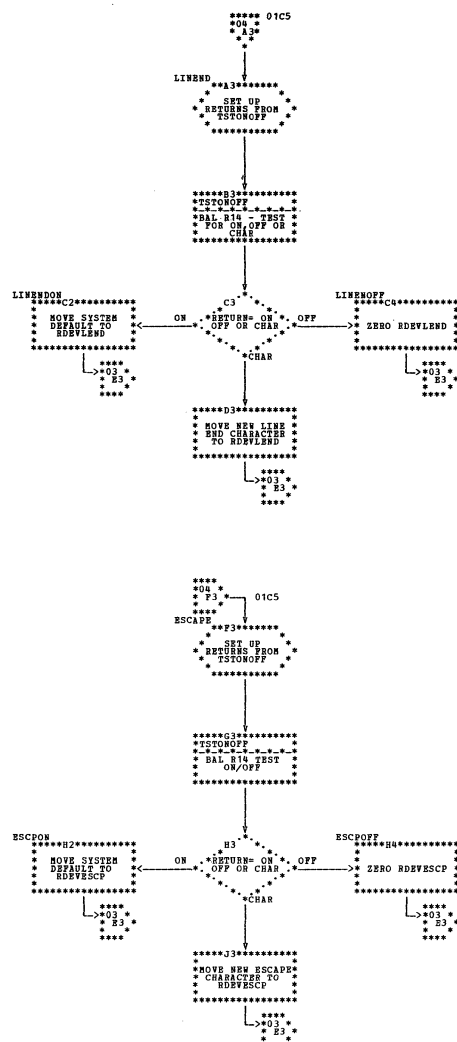
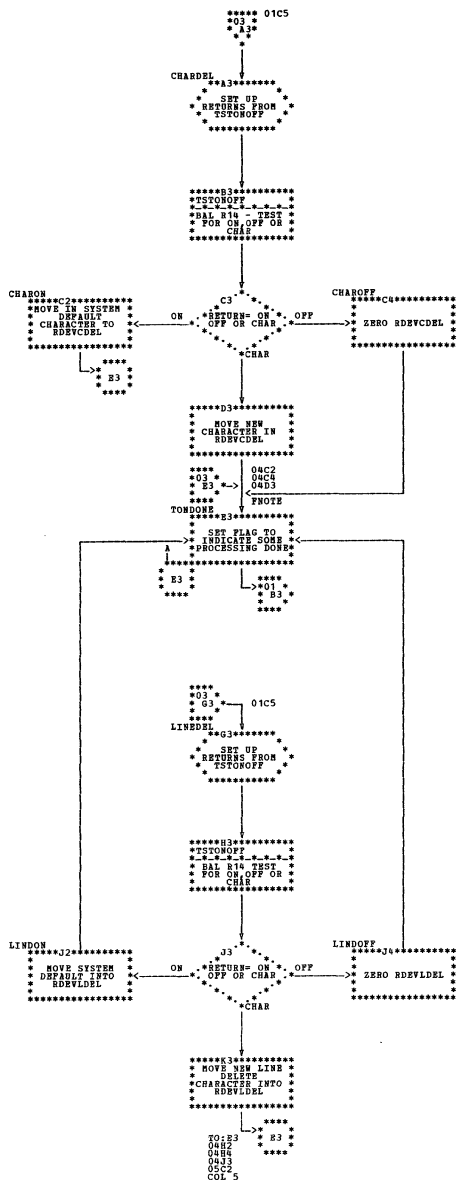






DMKCFM -- Process TERMINAL Command (Parts 1 and 2 of 6)

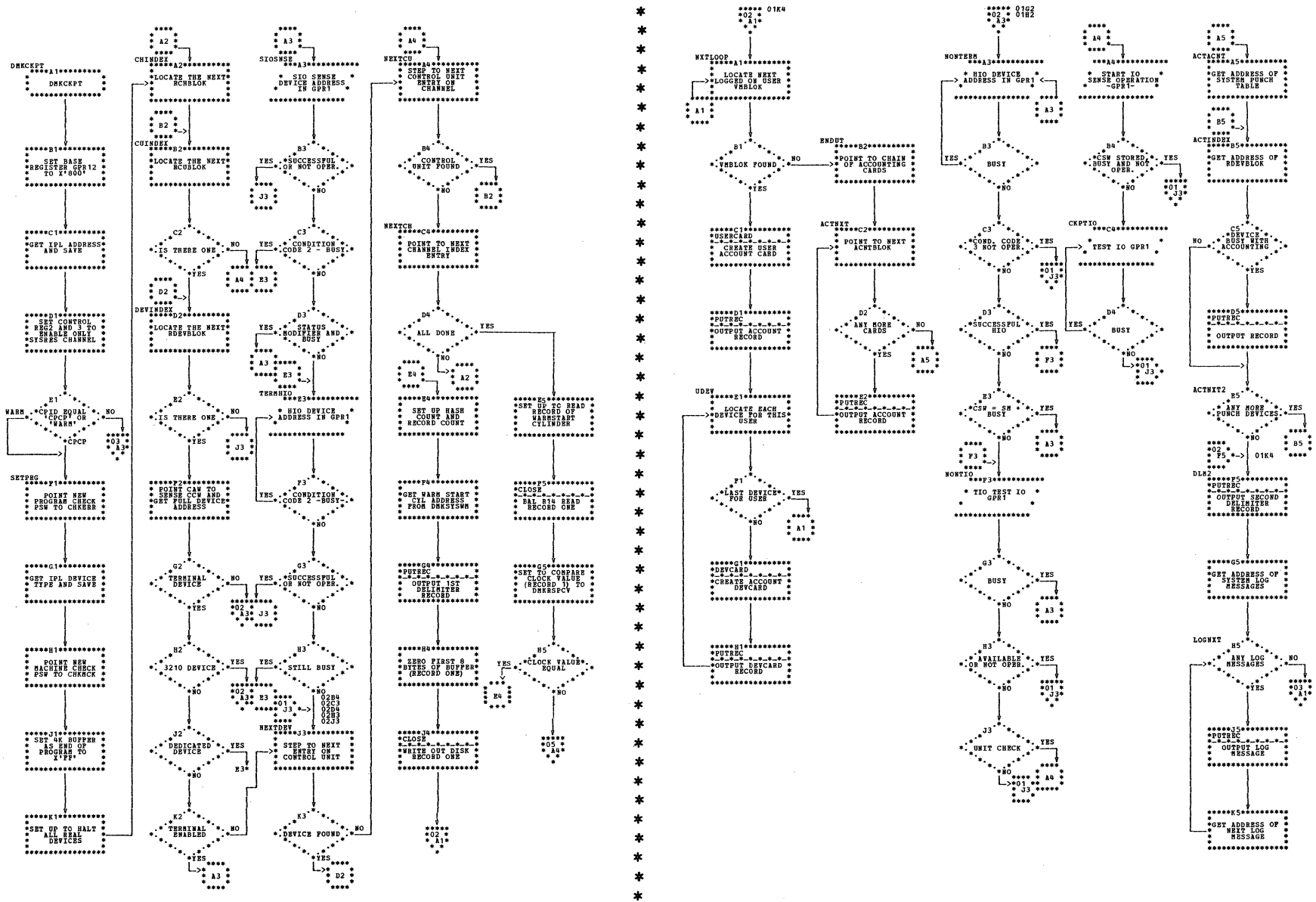
DMKCFT -- Process TERMINAL Command (Parts 3 and 4 of 6)

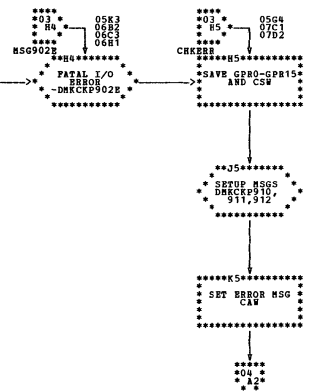
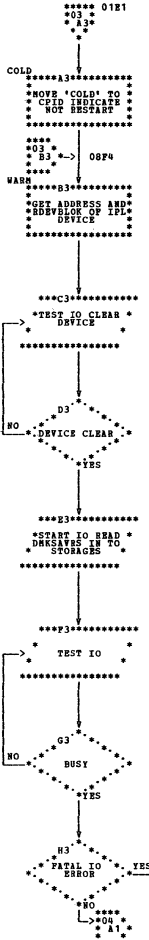
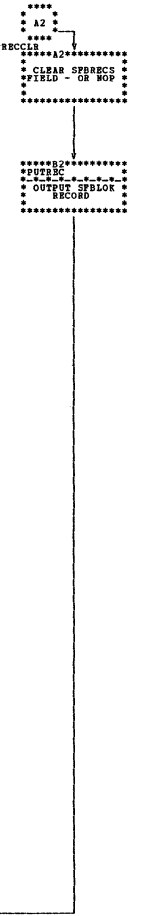
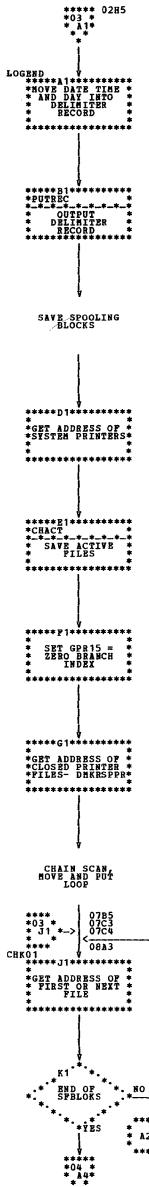


TO: E3  
04C4  
05C2  
05C4  
05E2  
05E4  
06C1  
06C3  
06C4

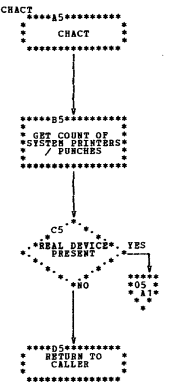
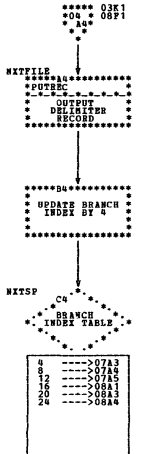
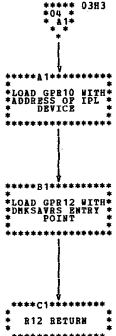


DMKCKP -- Checkpoint Processor (Parts 1 and 2 of 8)



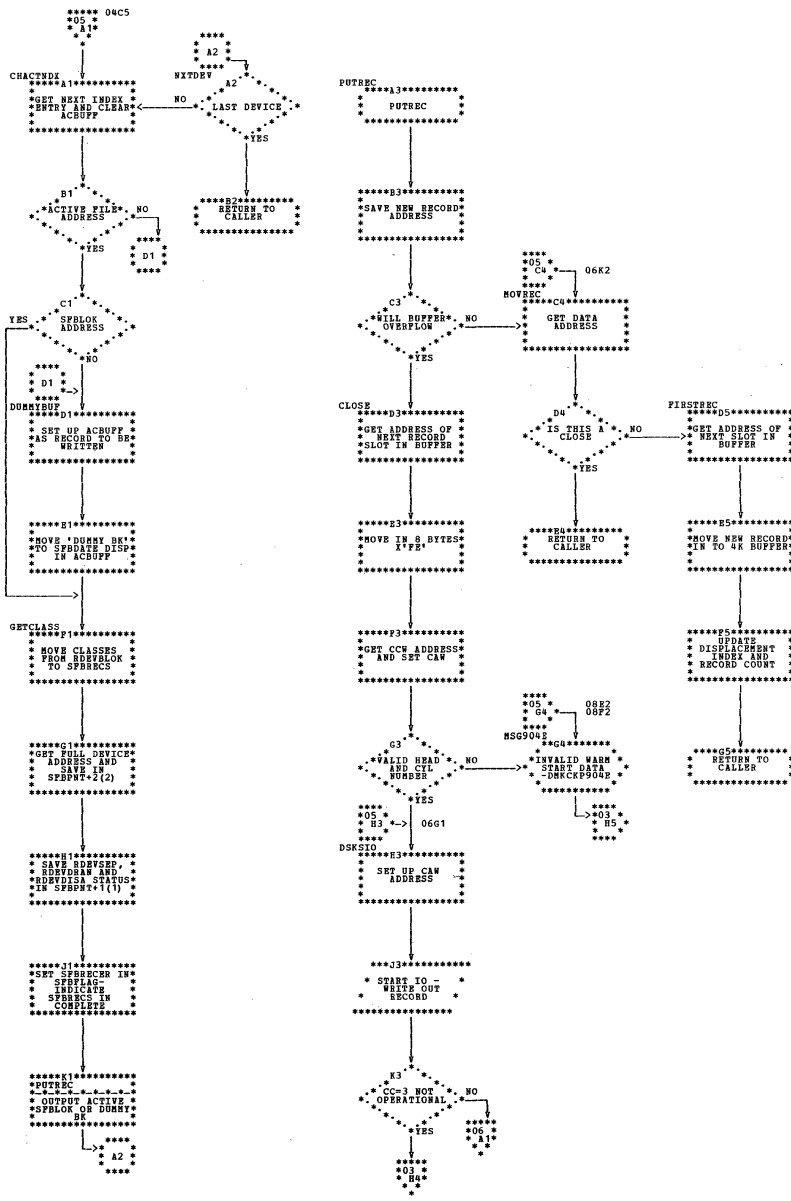


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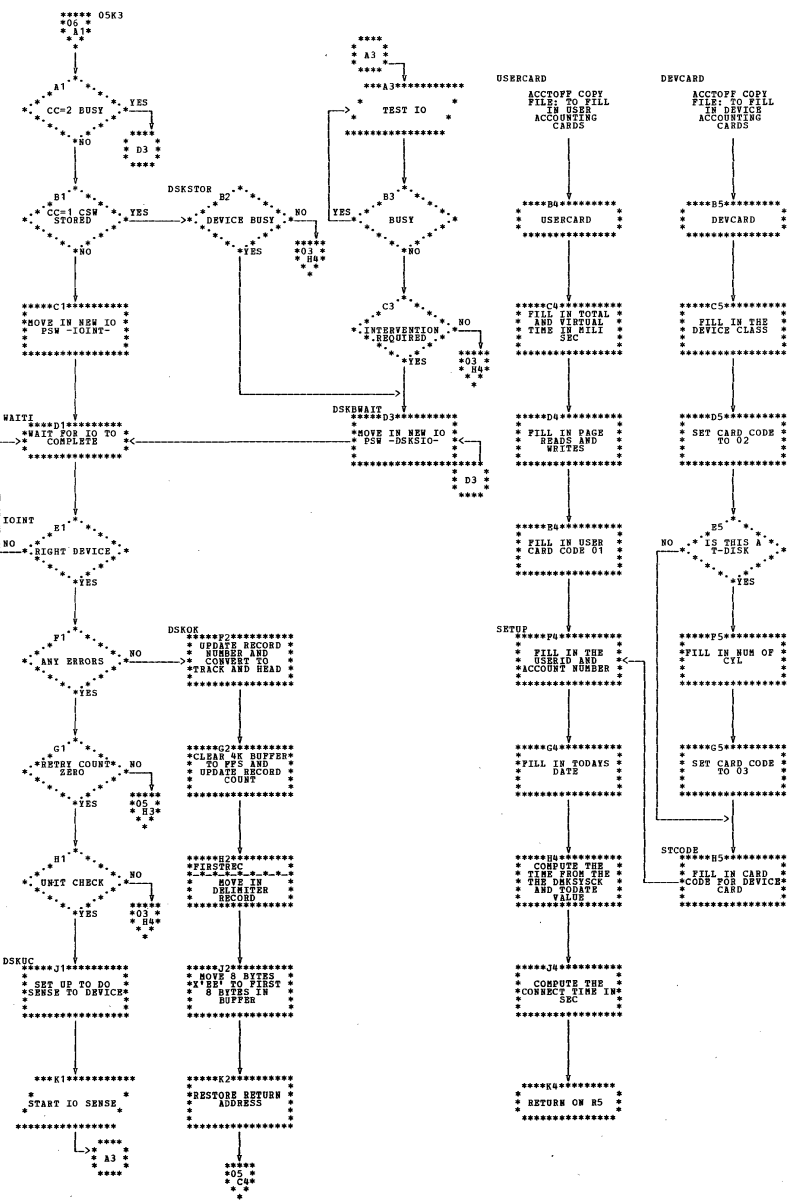


| DMKCKP -- Checkpoint Processor (Parts 3 and 4 of 8)

DMKCKP -- Checkpoint Processor (Parts 5 and 6 of 8)



\* \* \* \* \*











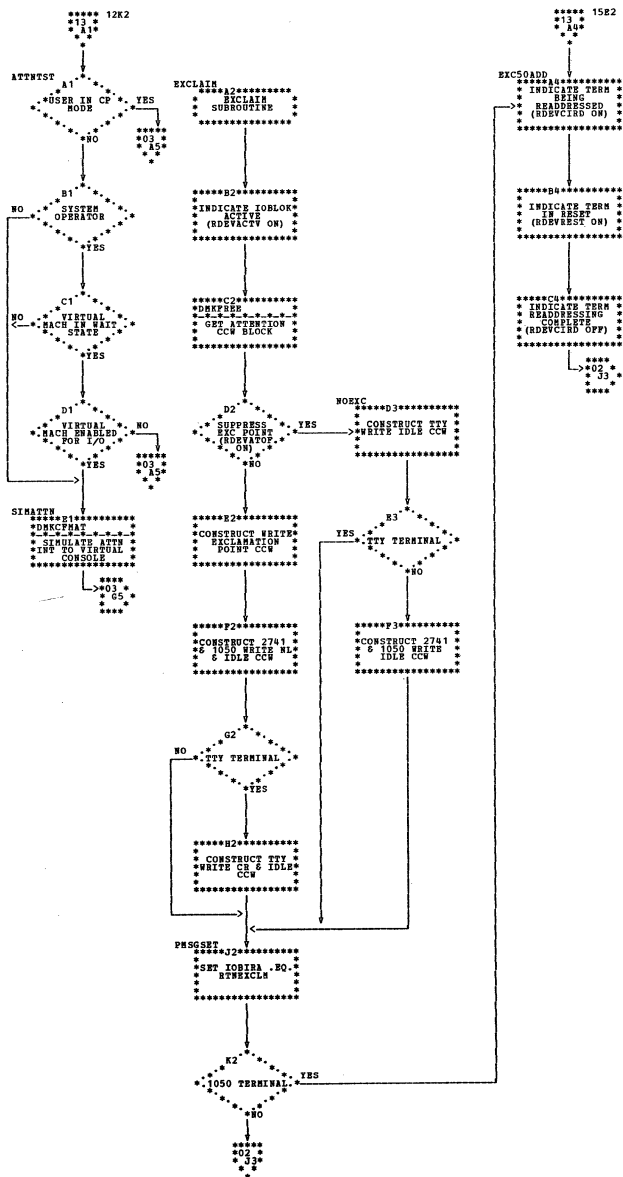




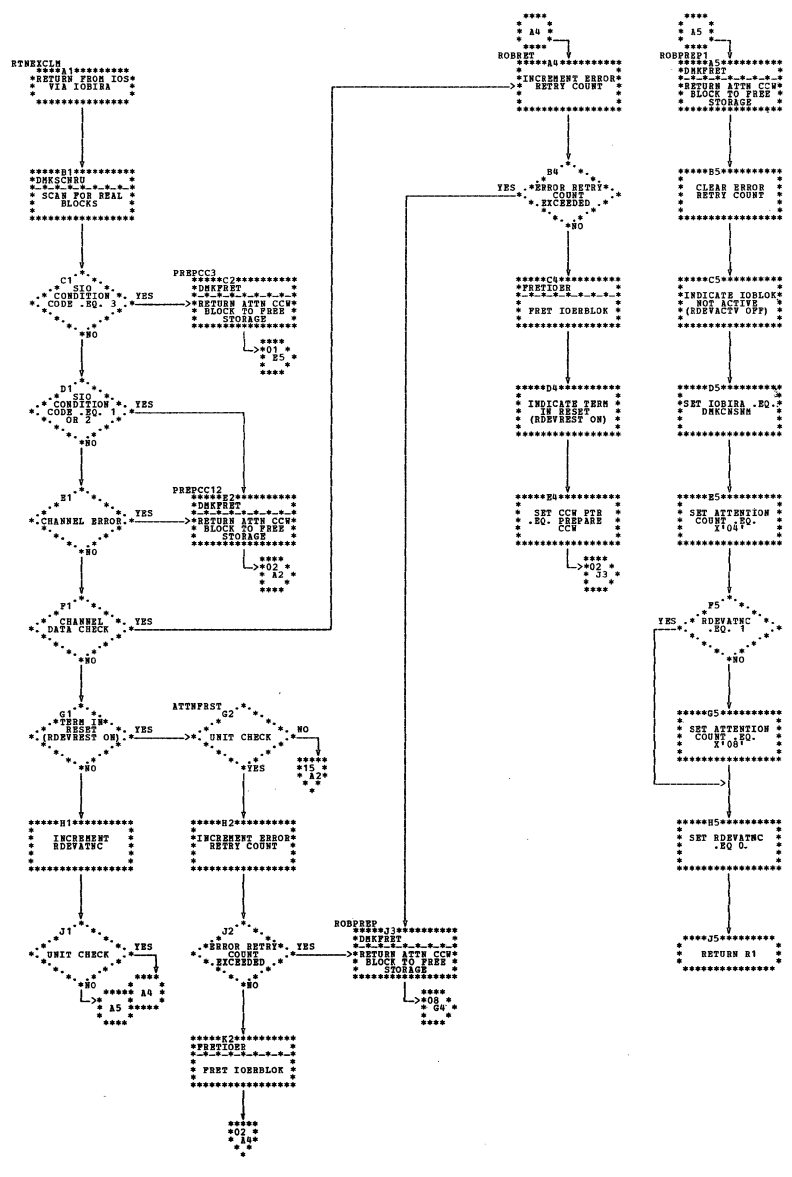




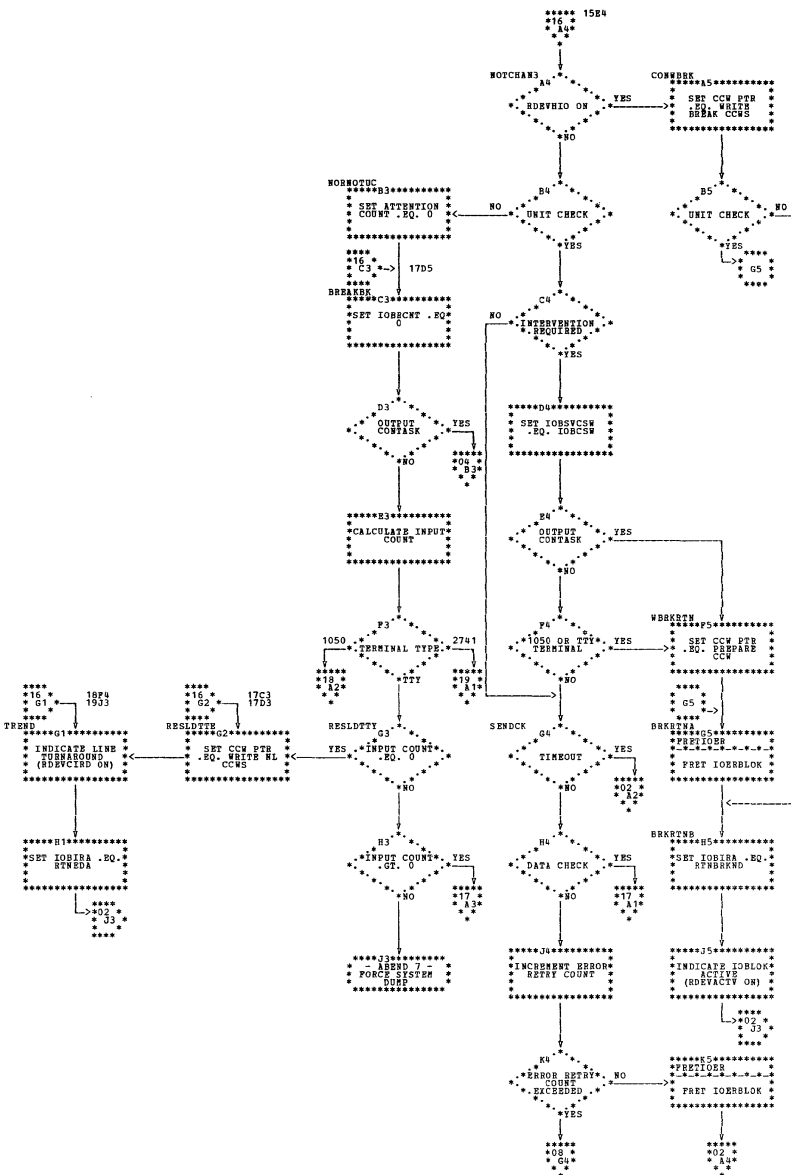
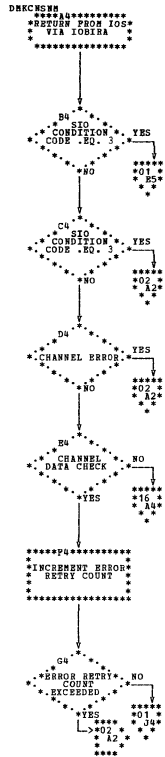
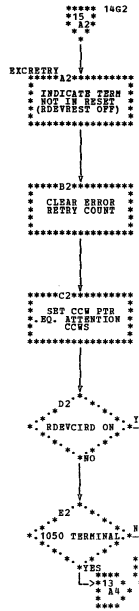
DMKONS -- Real Terminal (Console) Manager (Parts 13 and 14 of 23)



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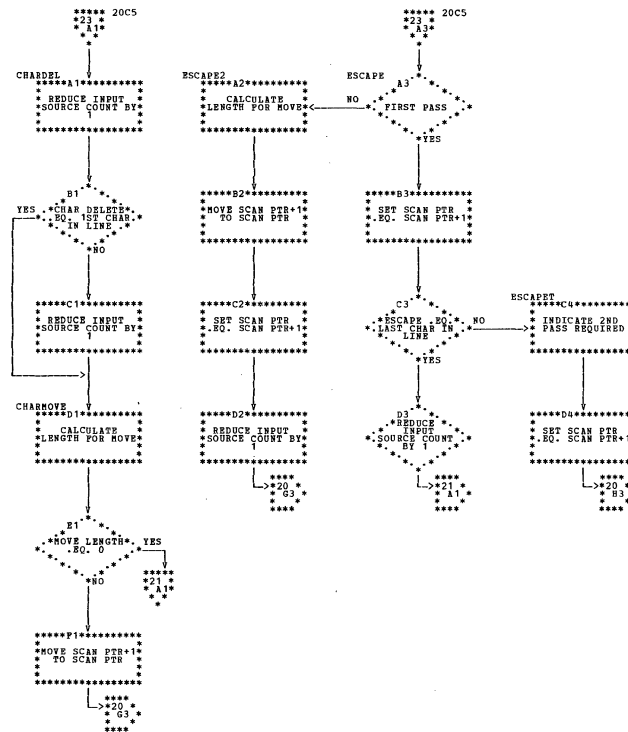




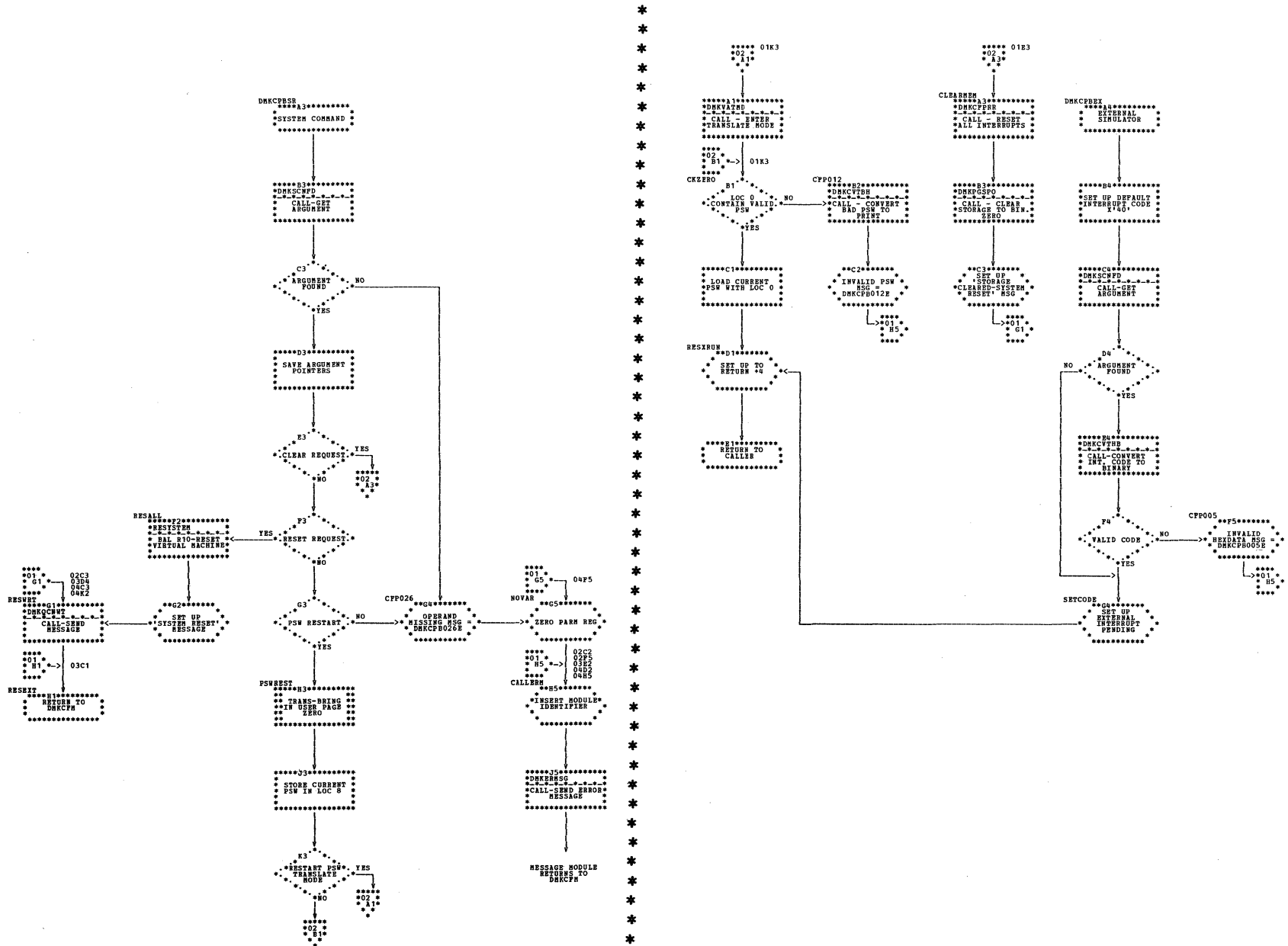


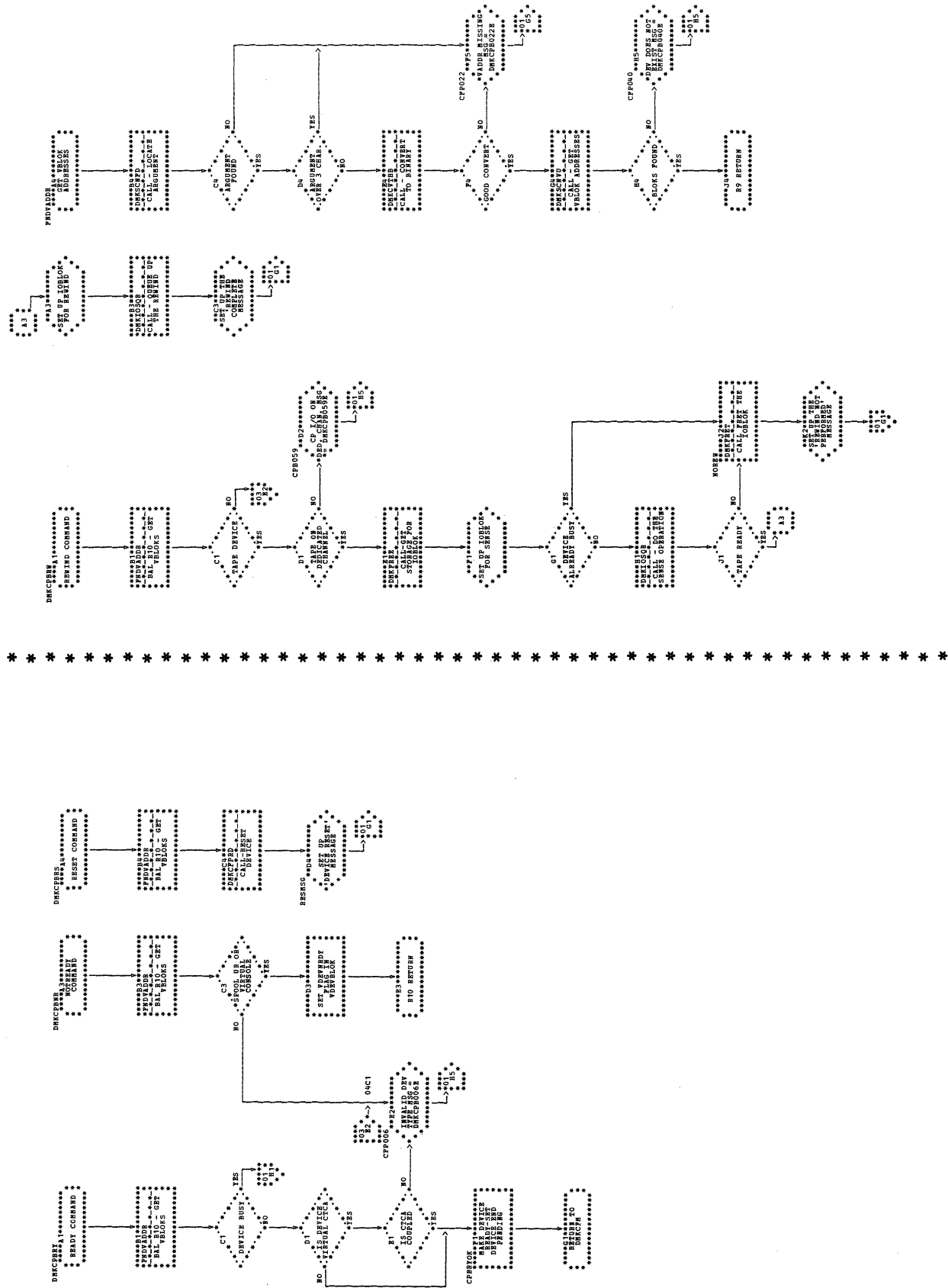






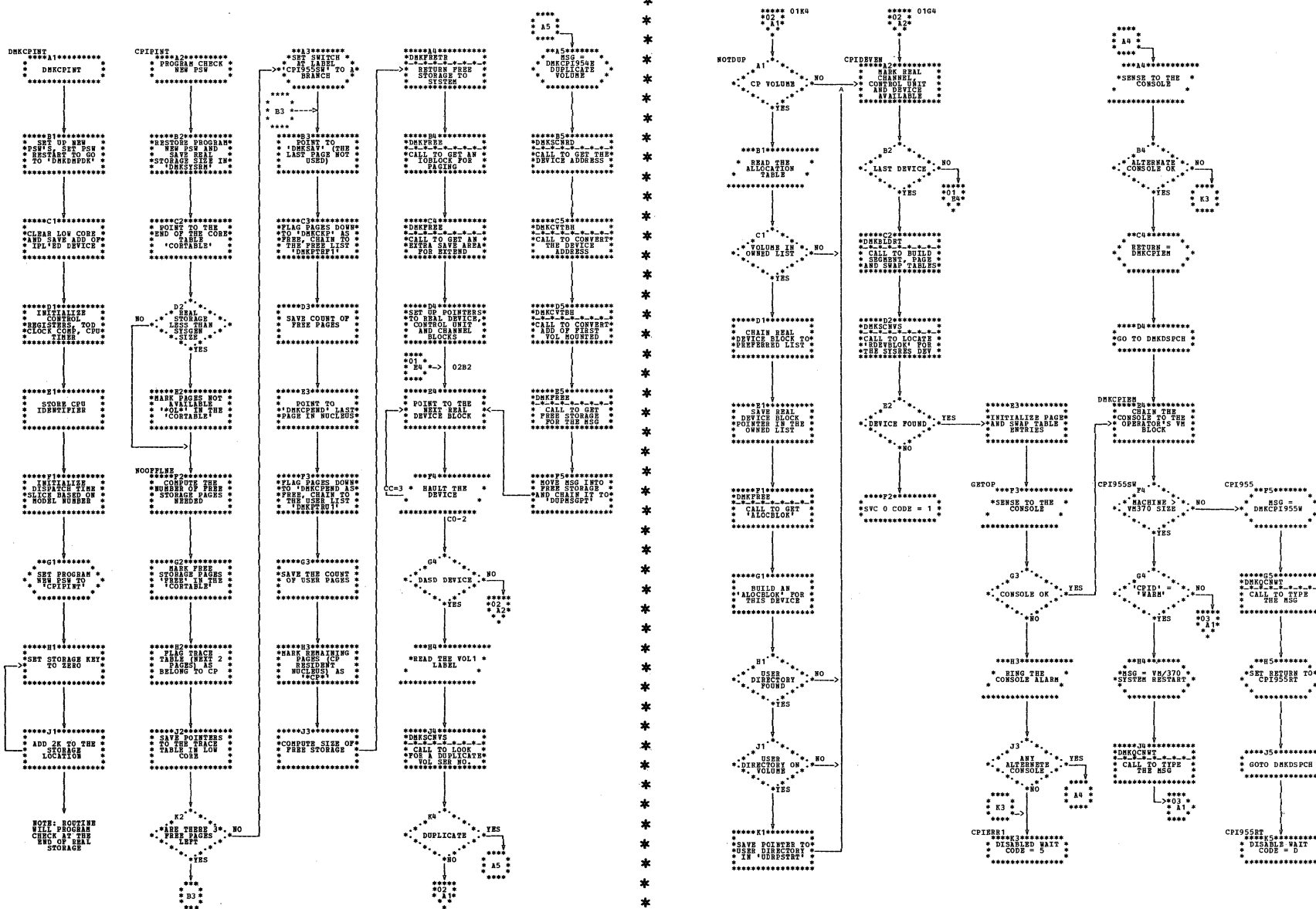
| DMKCPB -- Process EXTERNAL, NOTREADY, READY, REWIND, and SYSTEM Commands (Parts 1 and 2 of 4)



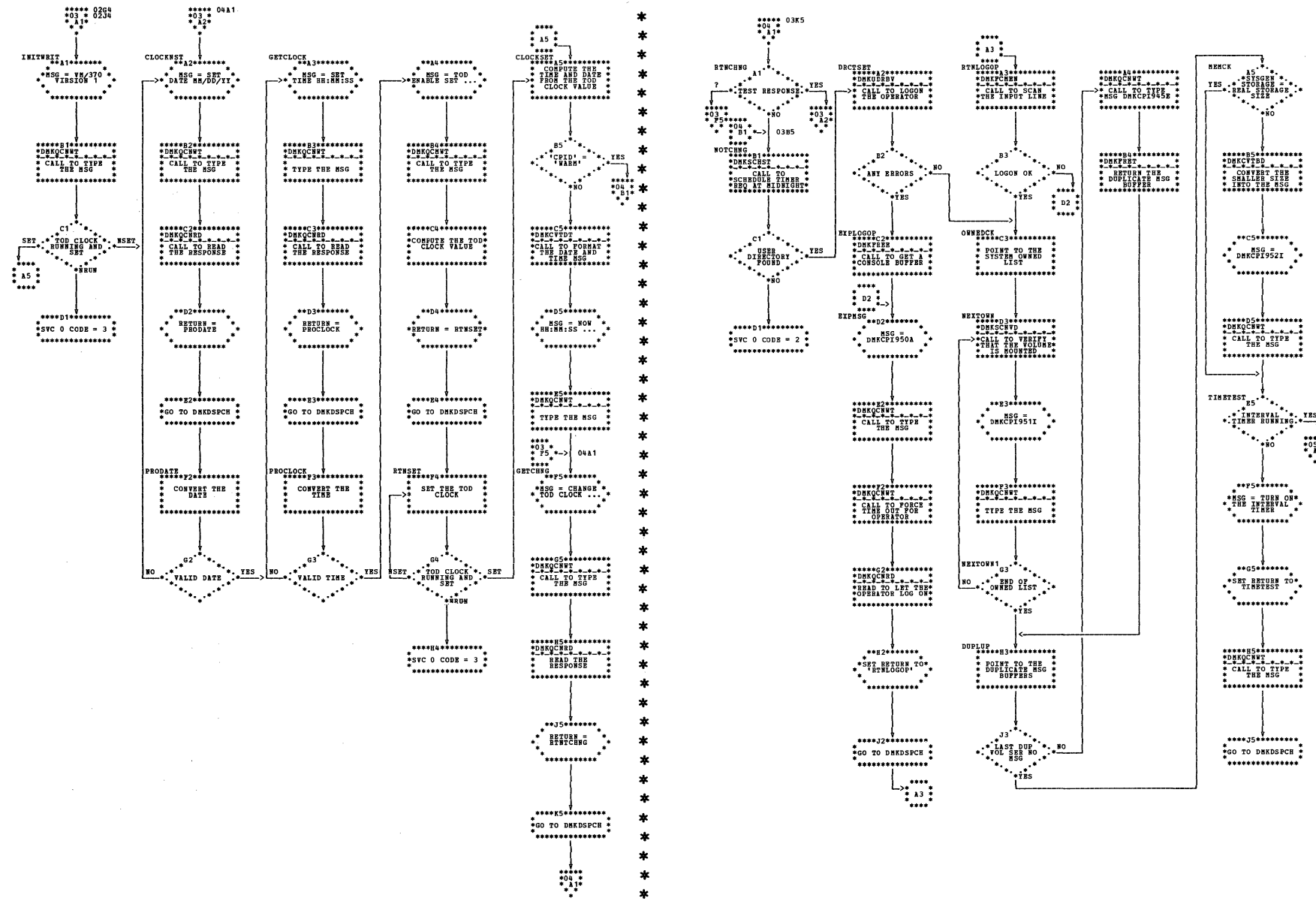


| DMKCPB -- Process EXTERNAL, NOTREADY, READY, REWIND, and SYSTEM COMMANDS (Parts 3 and 4 of 4)

DMKCPI -- Control Program Initialization (Parts 1 and 2 of 6)







| DMKCP1 -- Control Program Initialization (Parts 3 and 4 of 6)

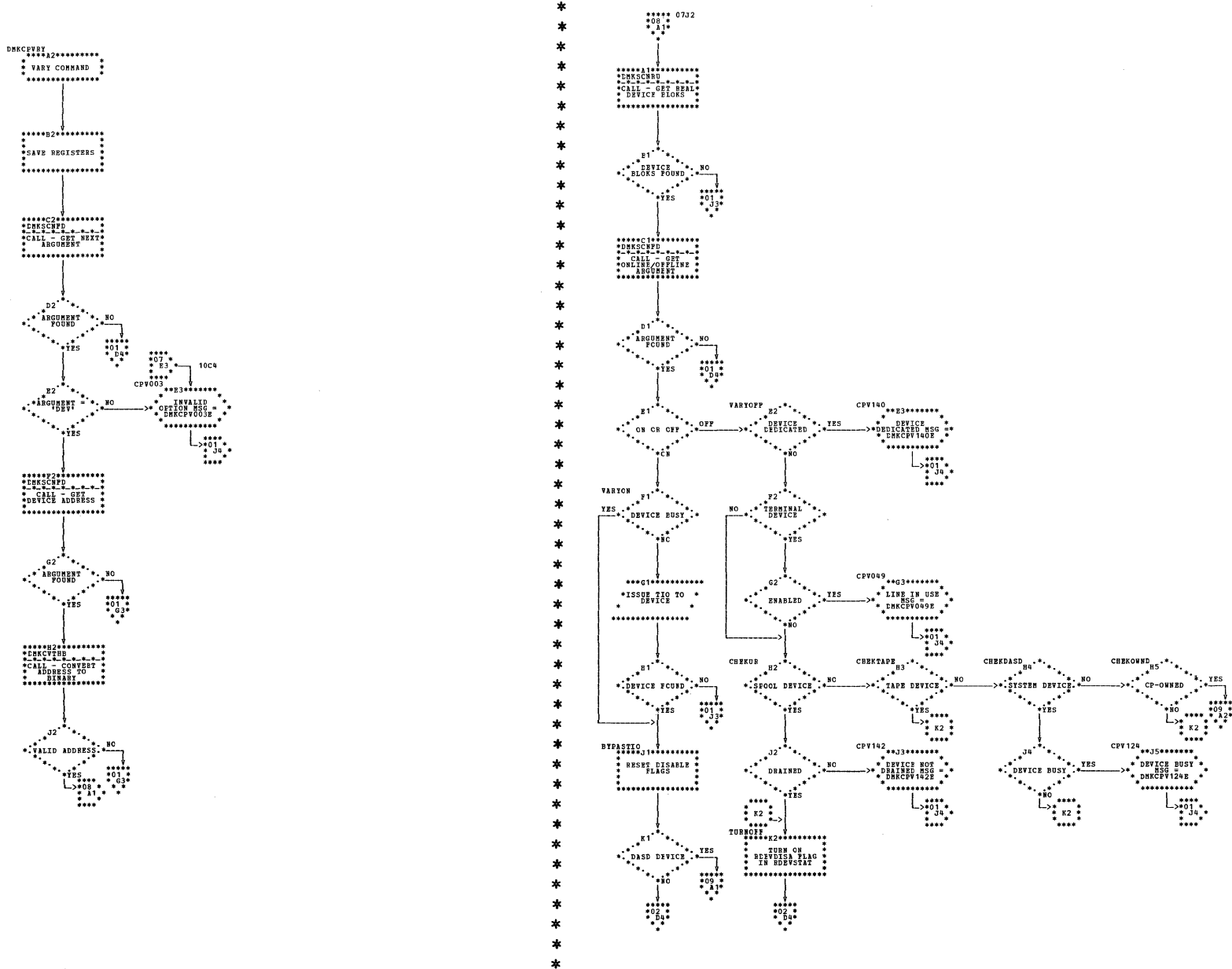


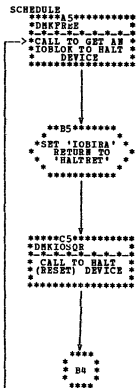
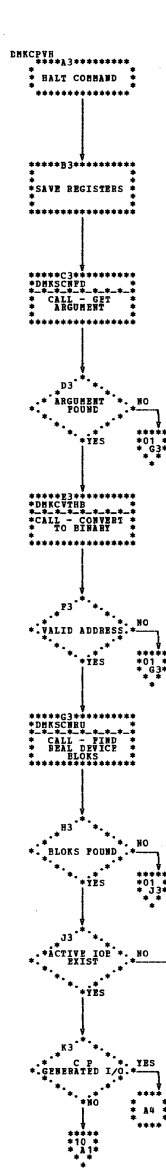
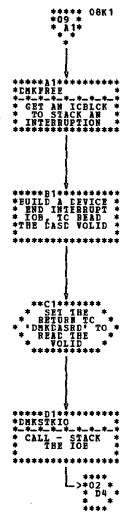




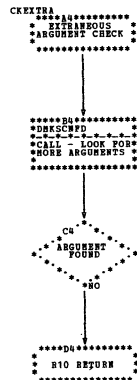
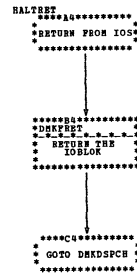
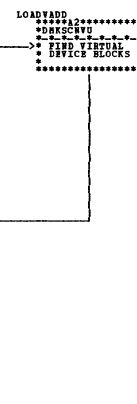


| DMKCPV -- Process DISABLE, ENABLE, HALT, SHUTDOWN, UNLOCK, and VARY Commands (Parts 7 and 8 of 10)





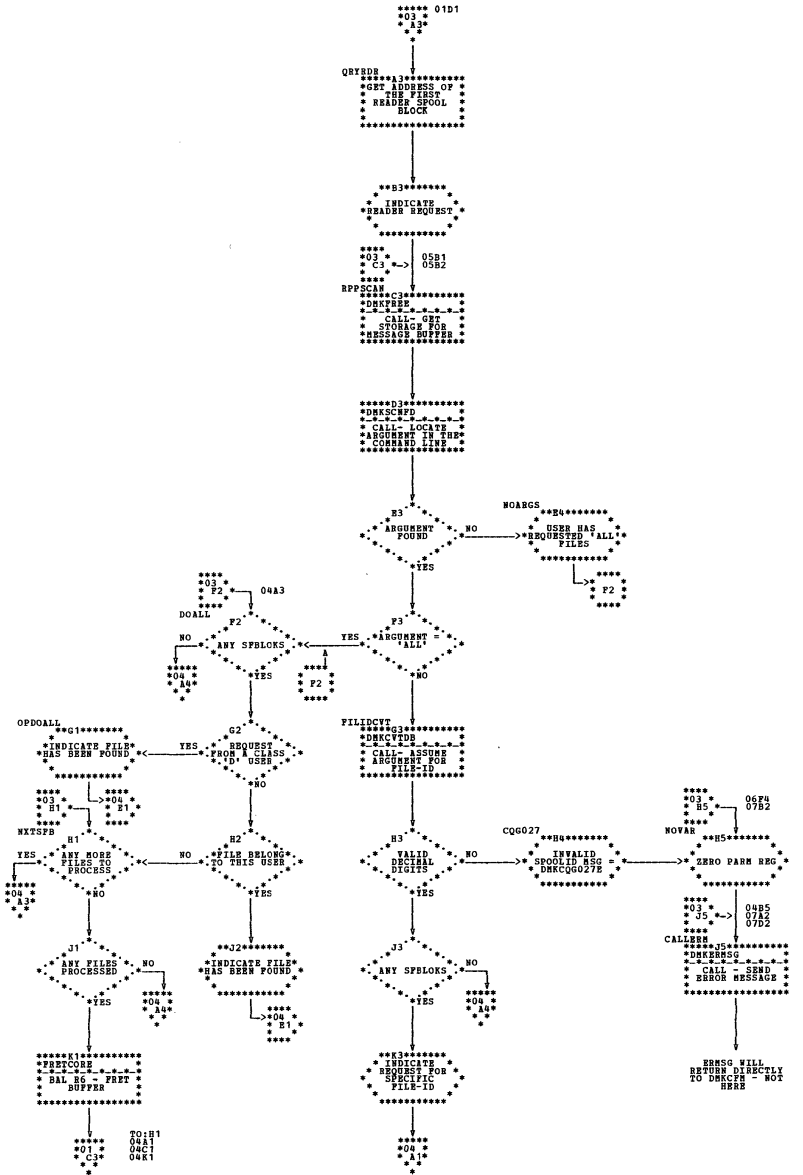
\* \* \* \* \*



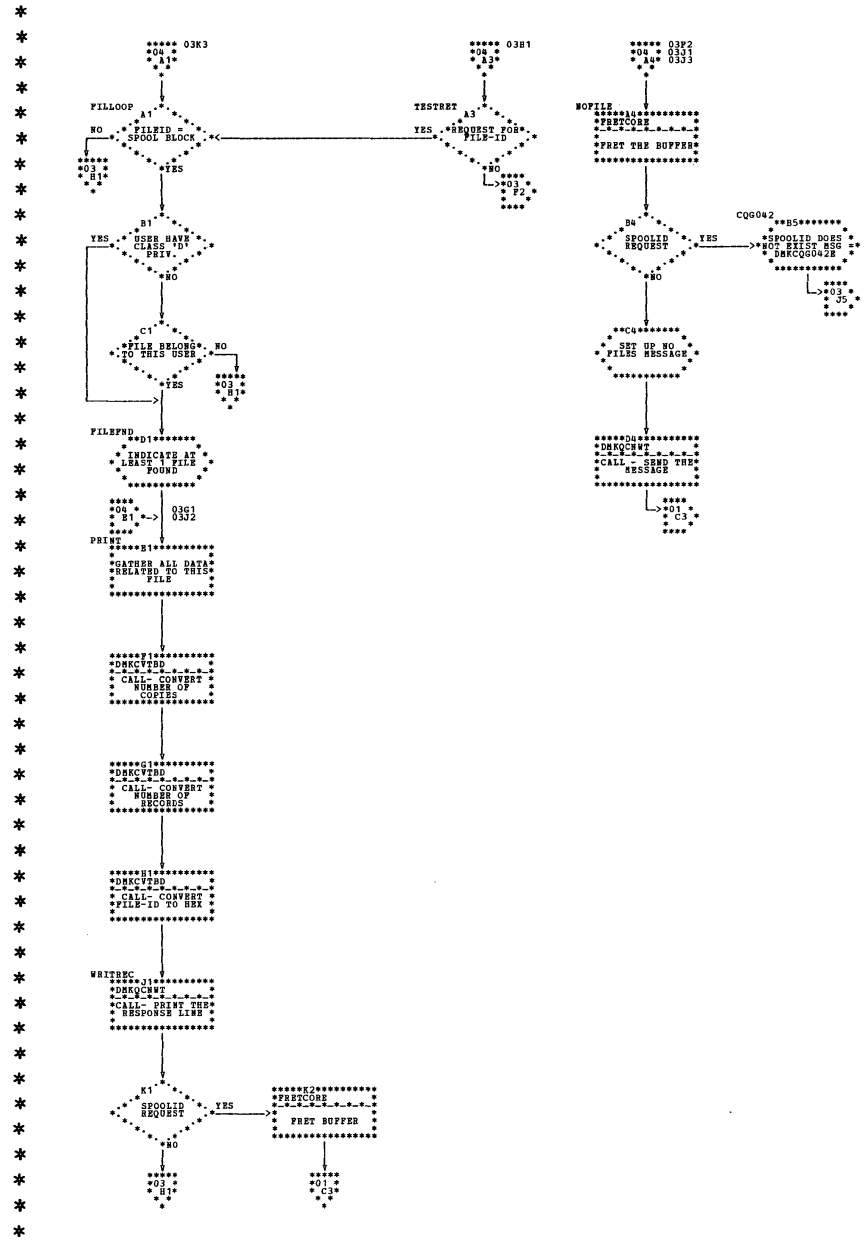
| DNKCPV -- Process DISABLE, ENABLE, HALT, SHUTDOWN, UNLOCK, and VARY Commands (Parts 9 and 10 of 10)



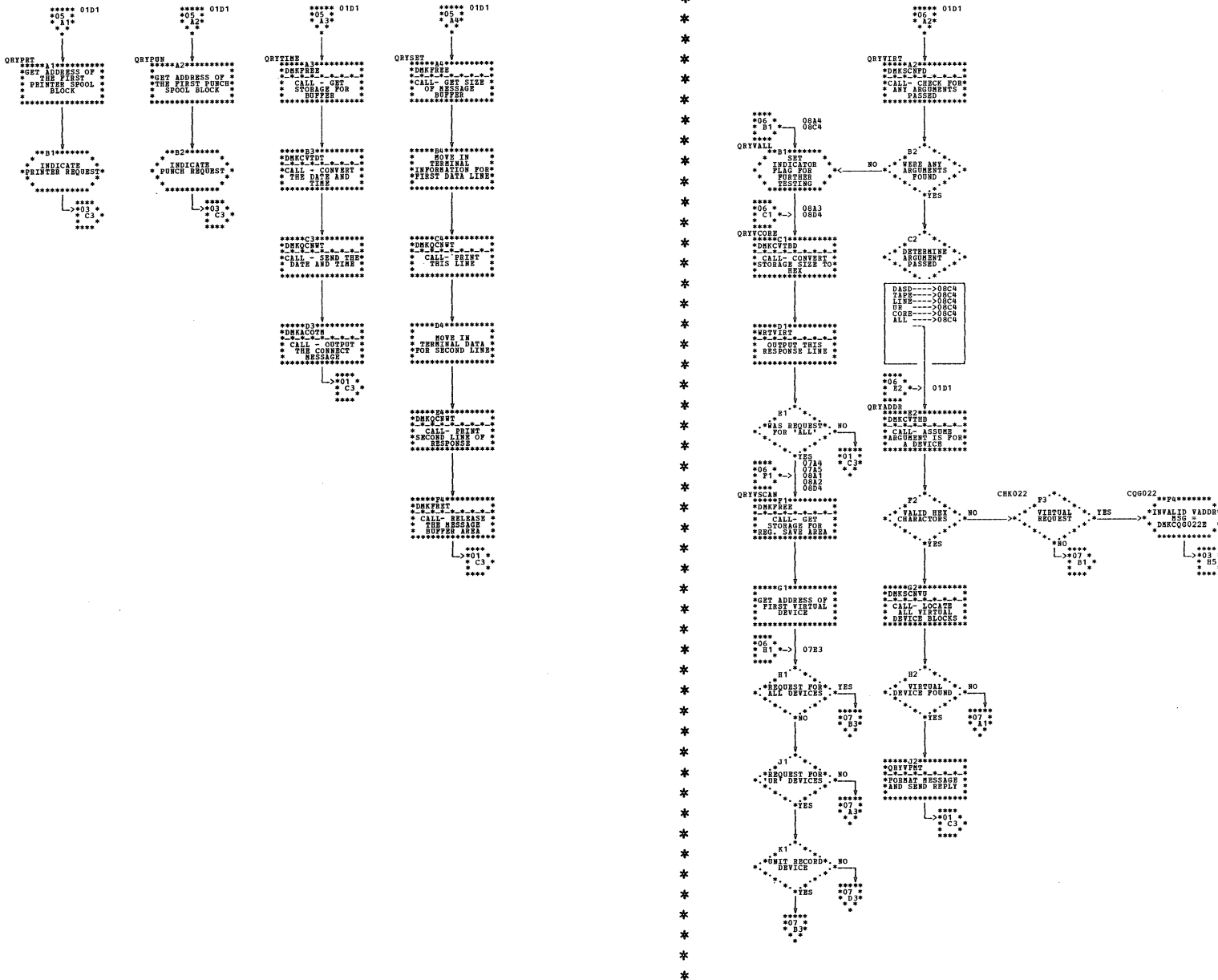




DMKCQG -- Process QUERY Command (Parts 3 and 4 of 10)

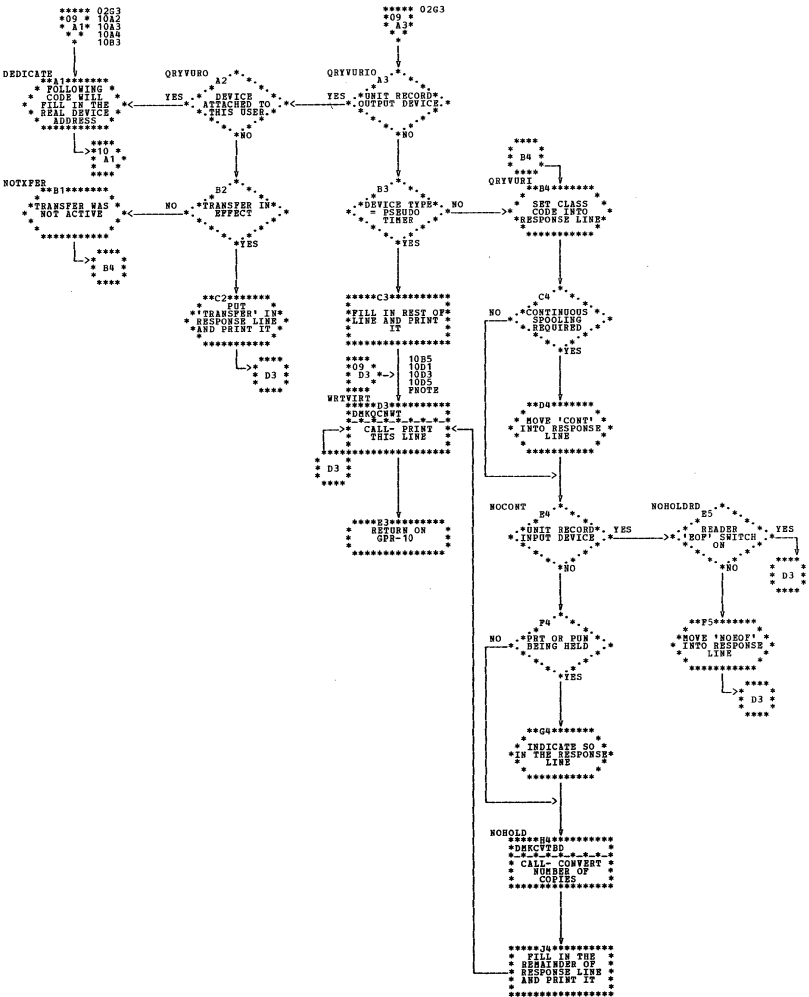


DMKCQG -- Process QUERY Command (Parts 5 and 6 of 10)

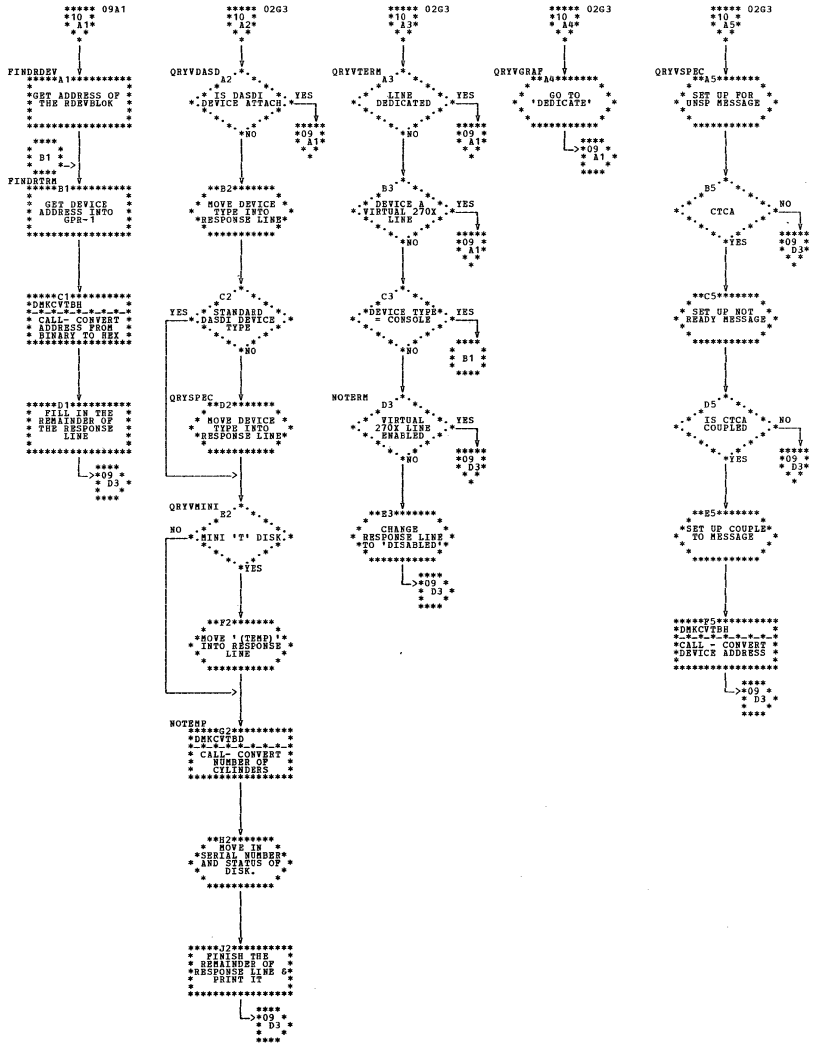




DMKQCQ -- Process QUERY Command (Parts 9 and 10 of 10)



TO: D3  
1082  
10A3  
10B3



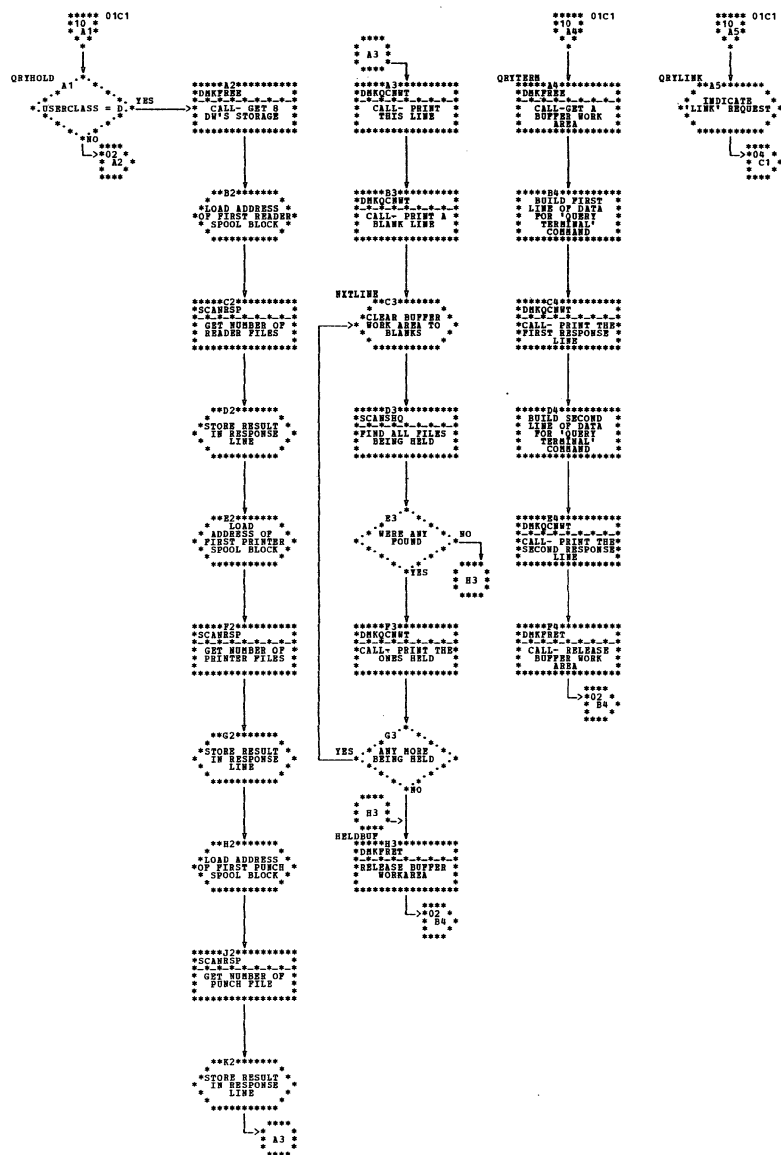
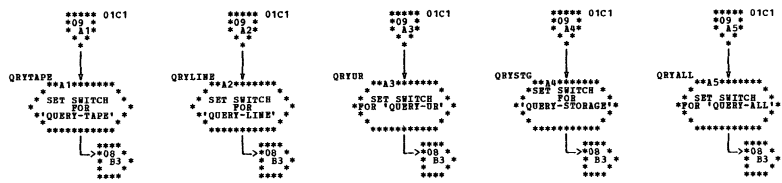






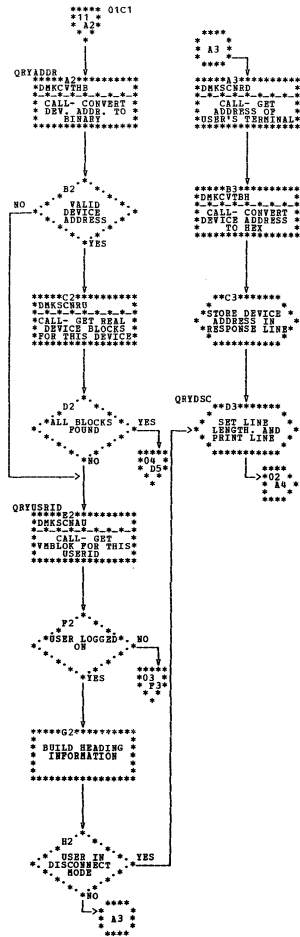


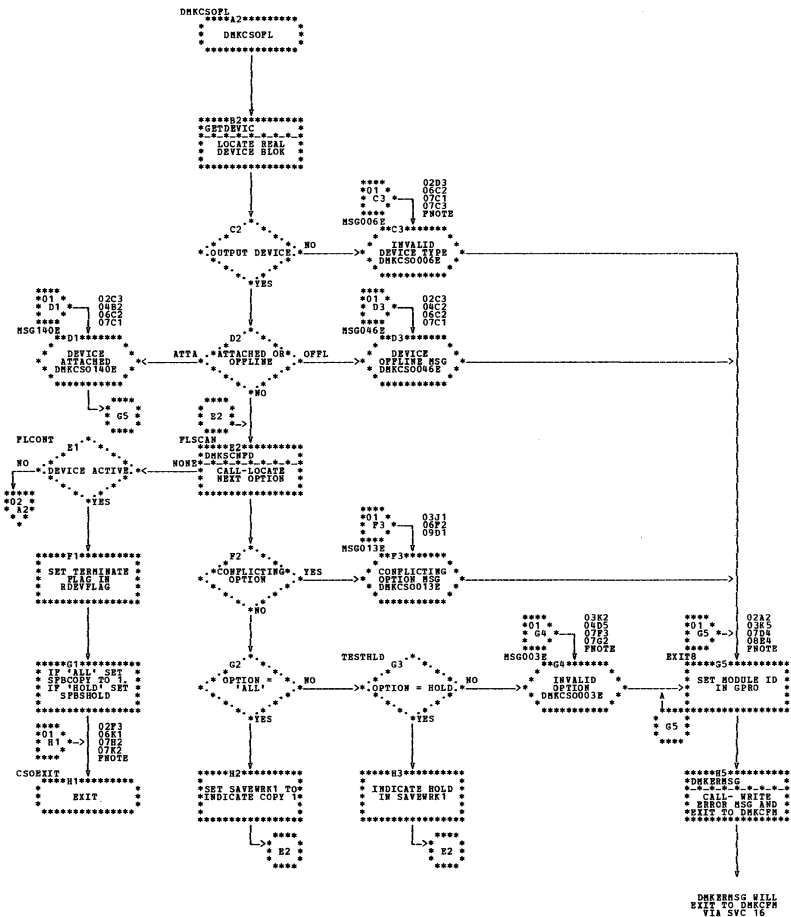




DMKCQP -- Process QUERY Command (Parts 9 and 10 of 11)

DMKCQP -- Process QUERY Command (Part 11 of 11)





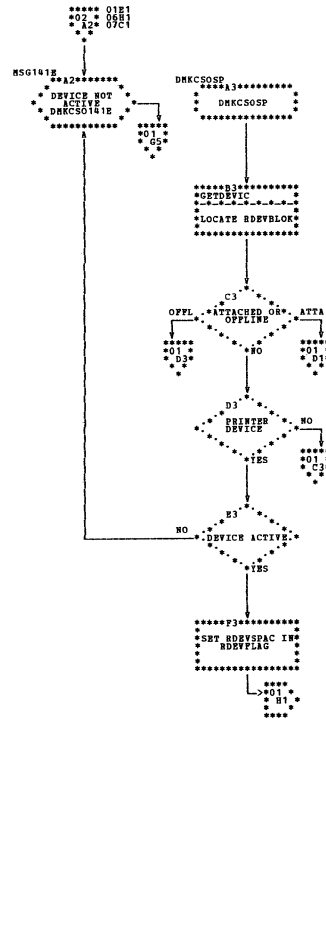
TO:G1  
1281  
1301  
1384

TO:G2  
1283  
1293

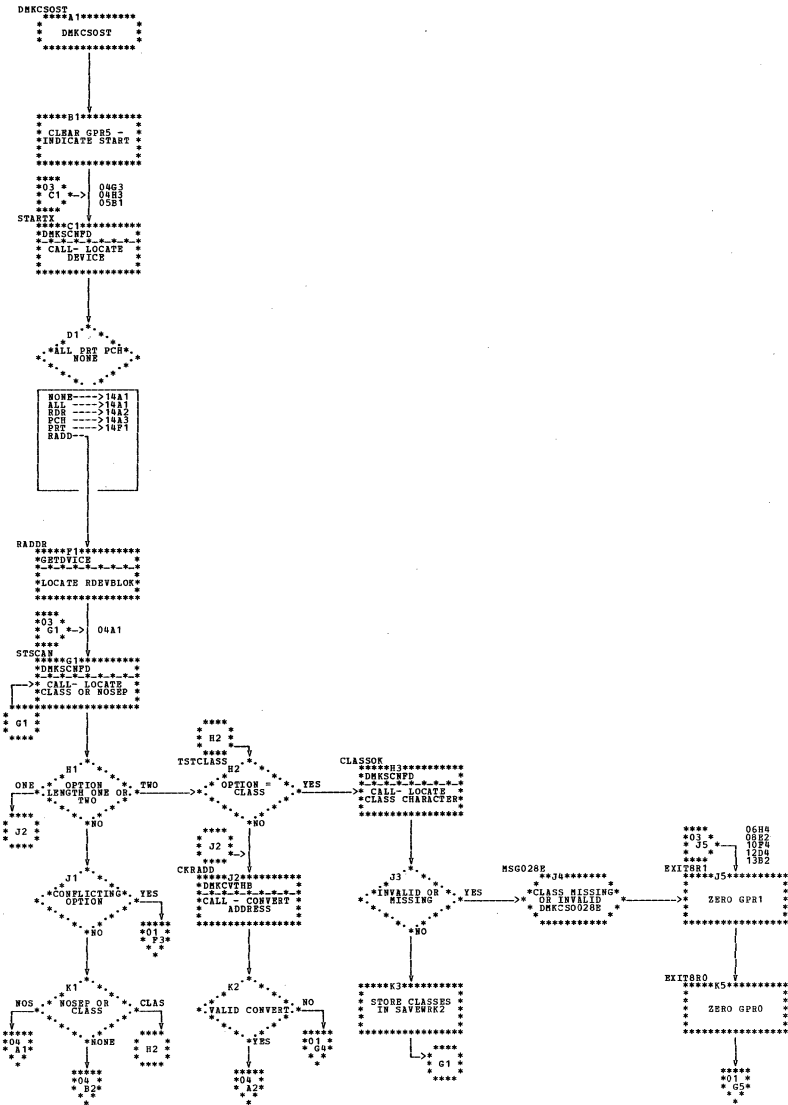
TO:G4  
1281  
1283  
1293

TO:G5  
1282  
1282

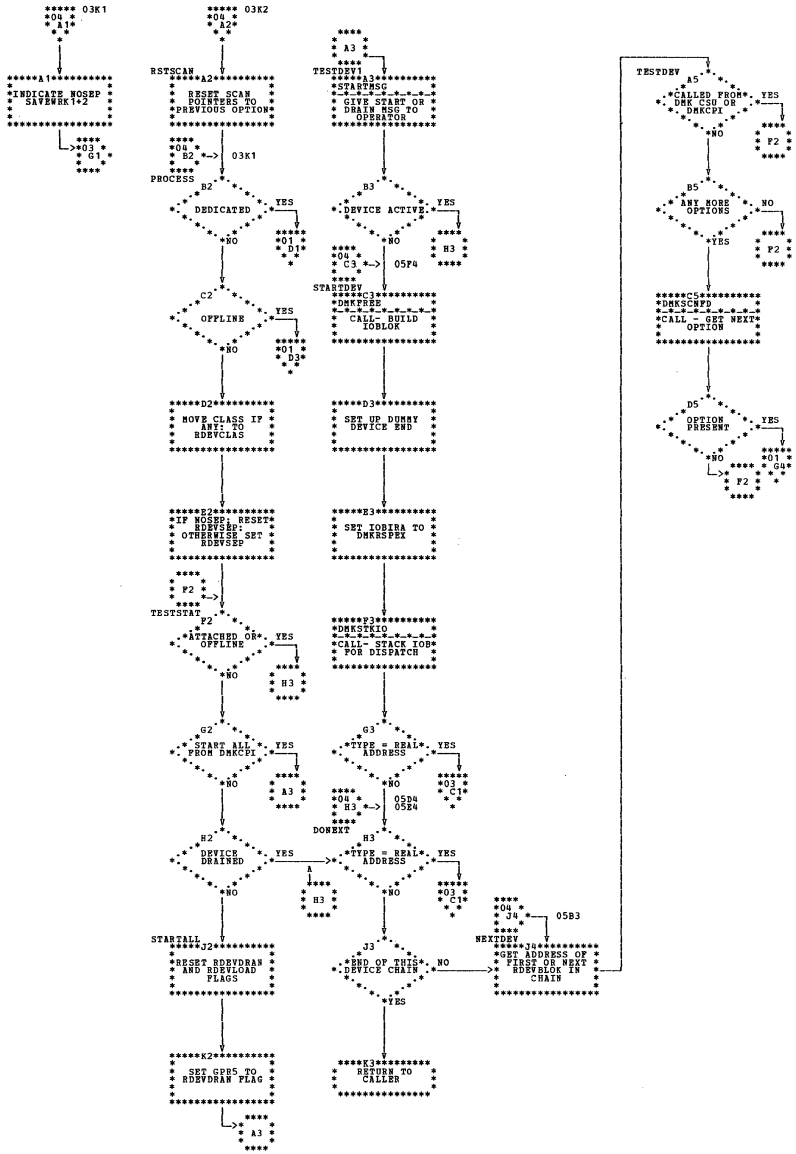
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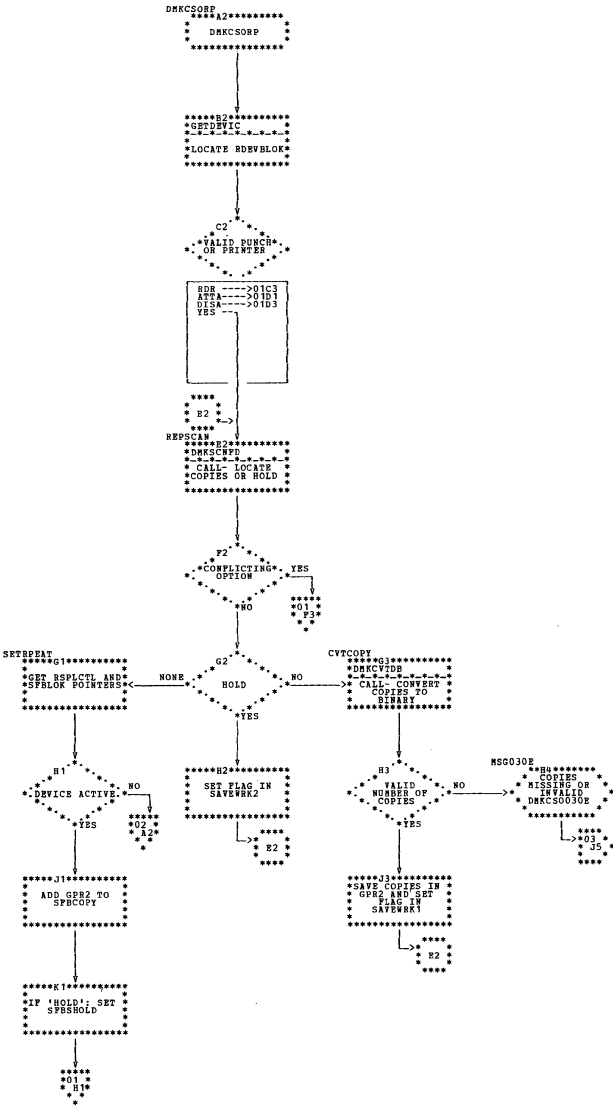
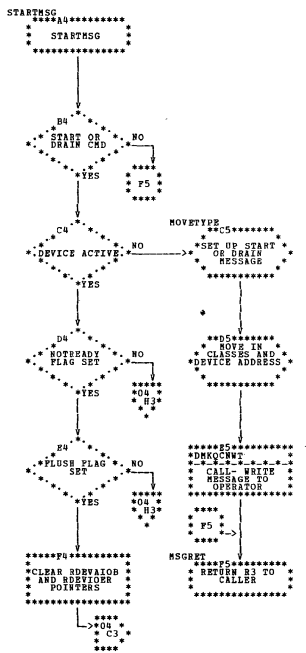
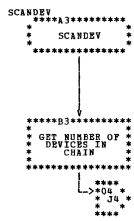
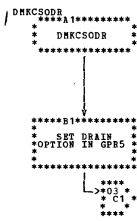


| DMKCSO -- Process Spooling Commands (Parts 3 and 4 of 14)



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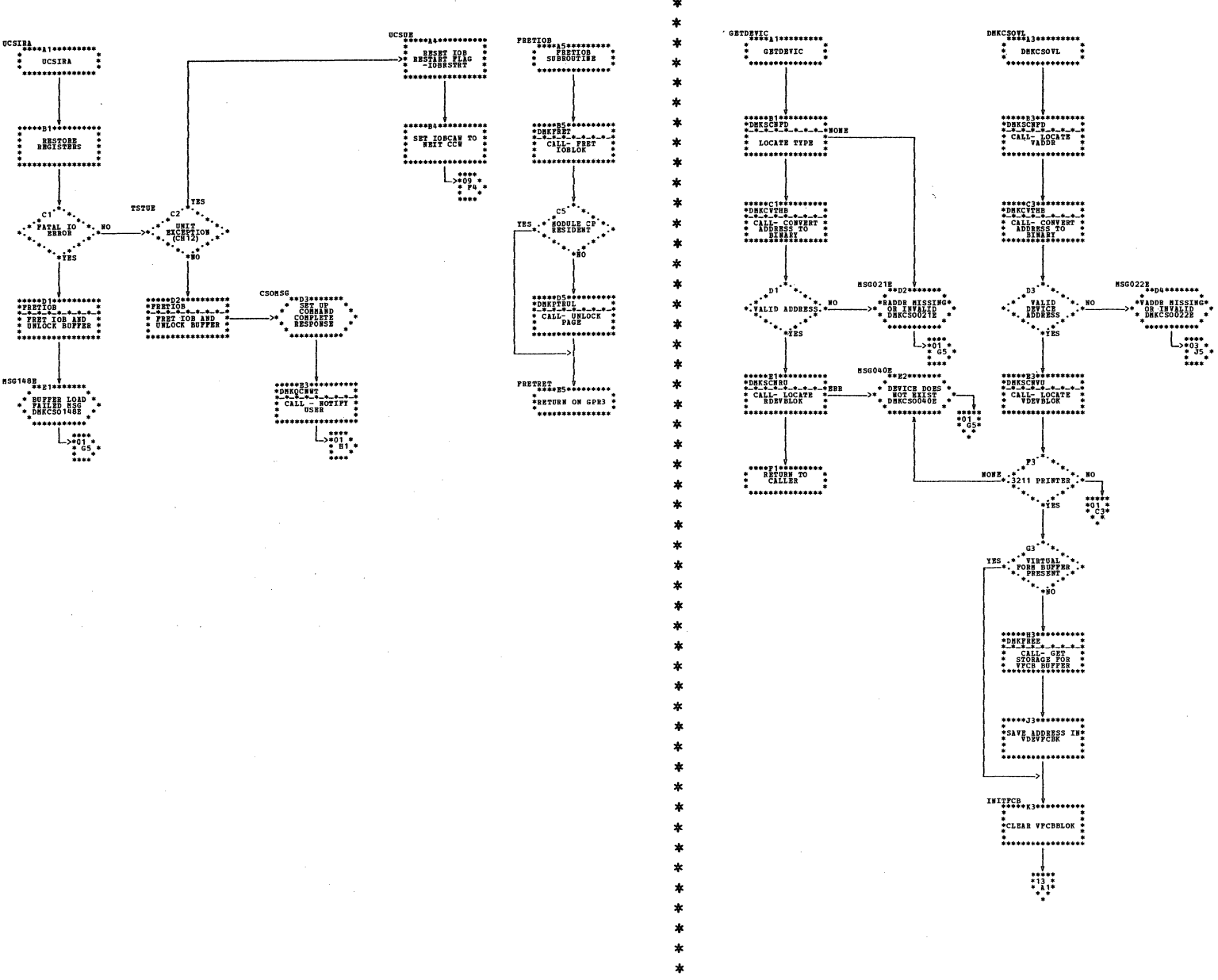




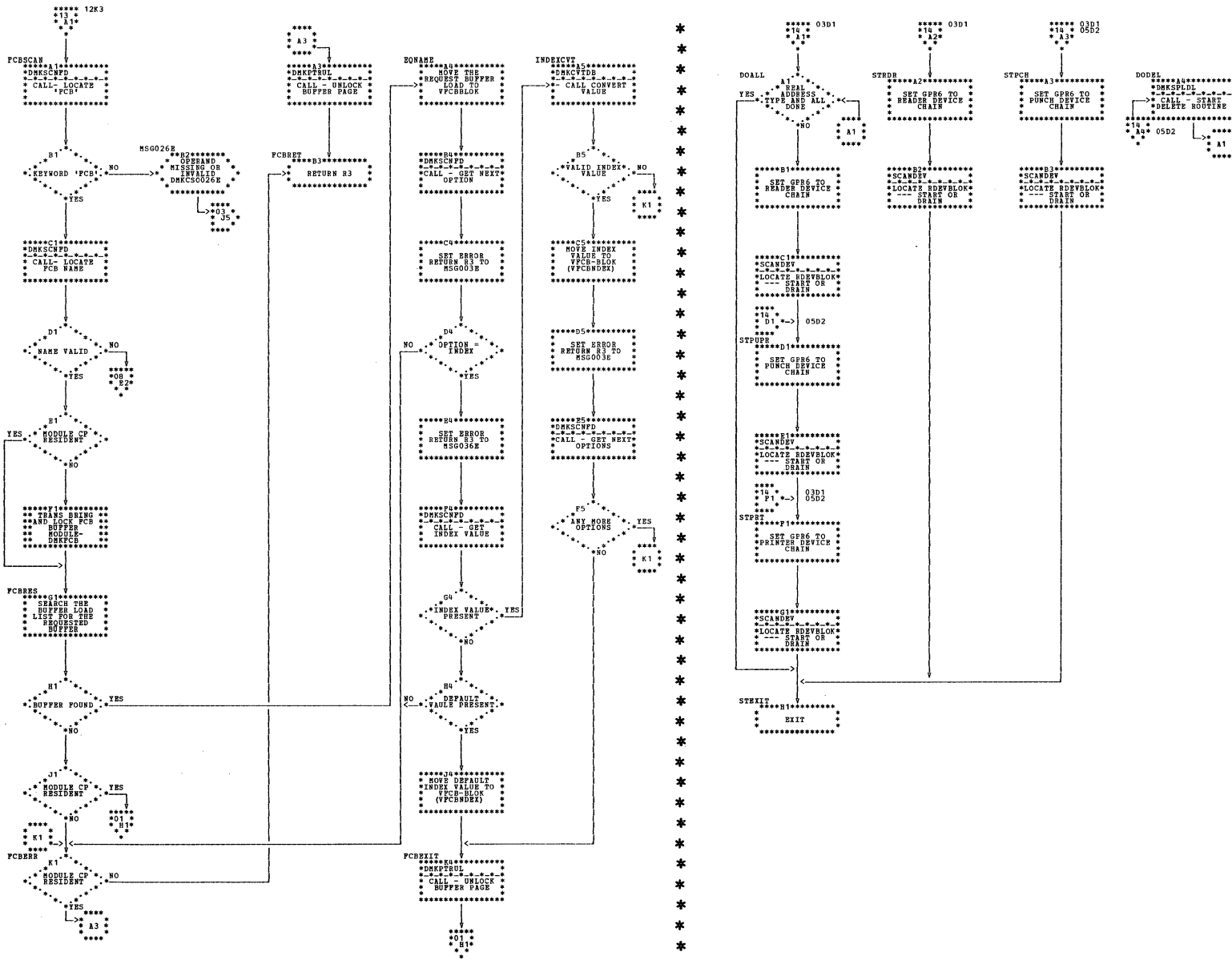




DMKCSC -- Process Spooling Commands (Parts 11 and 12 of 14)

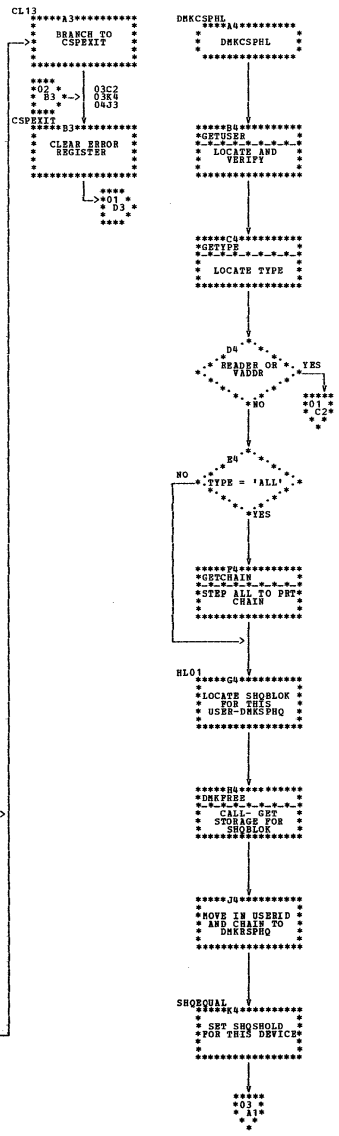
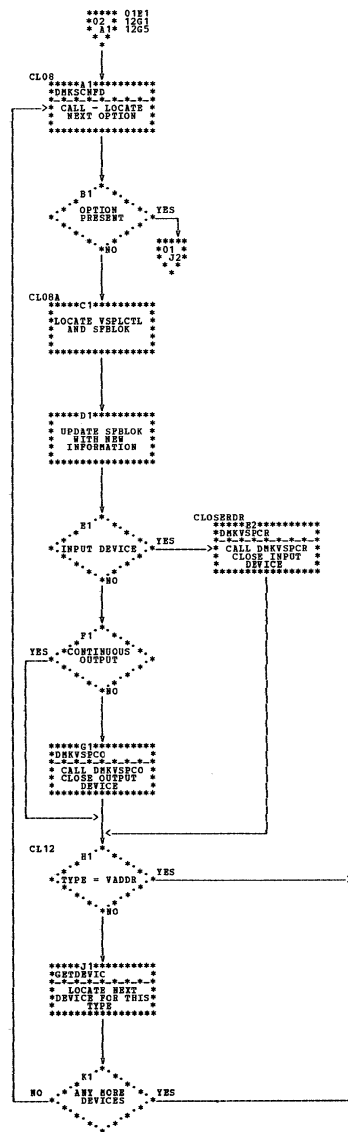
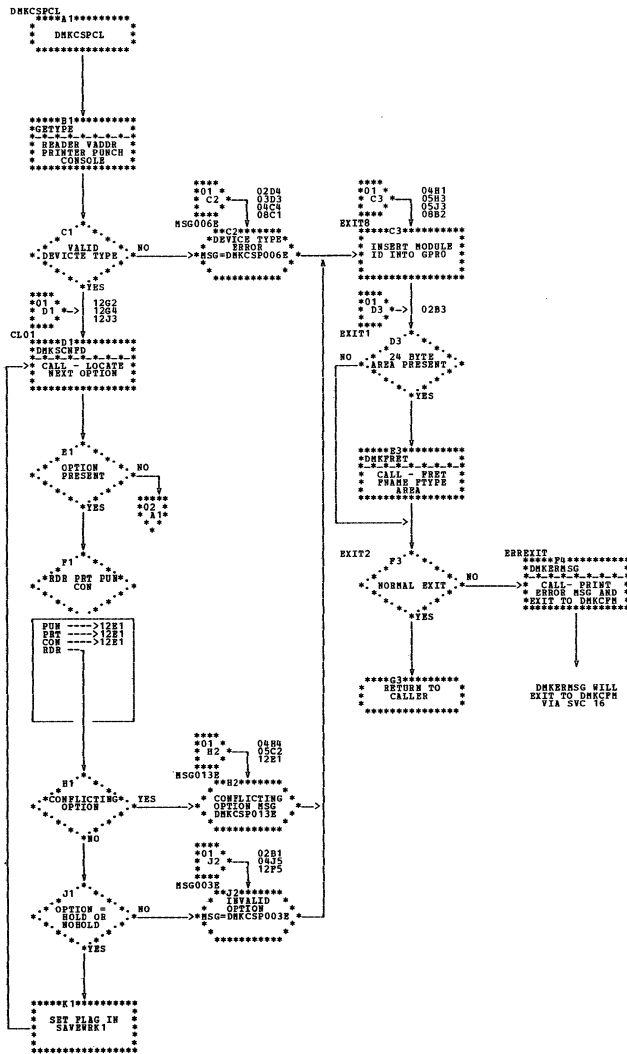


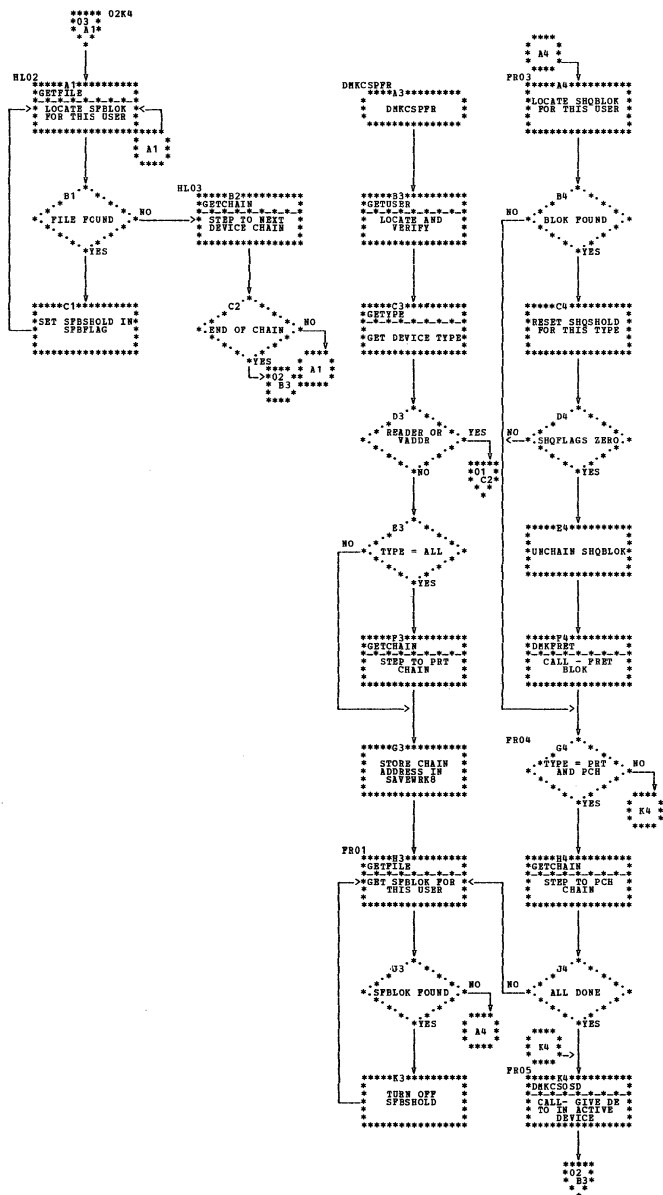




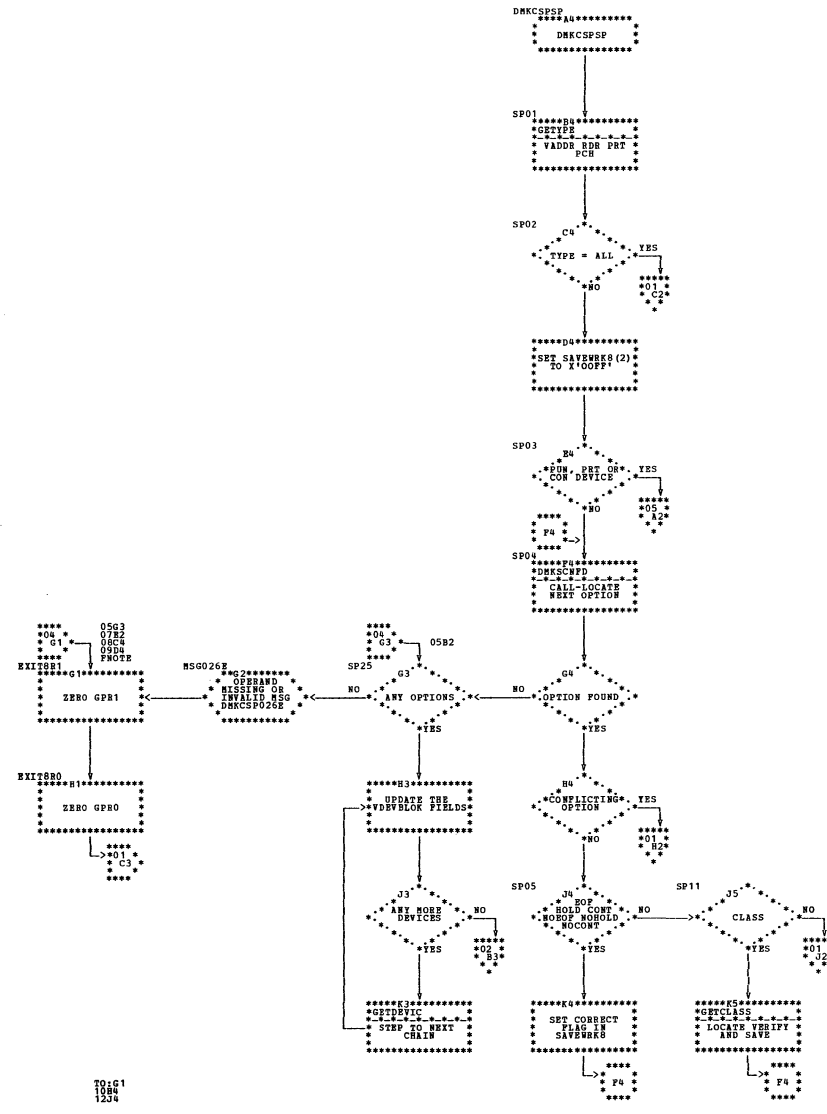
DMKCSO -- Process Spooling Commands (Parts 13 and 14 of 14)

DMKCSPL -- Process Spooling Commands (Parts 1 and 2 of 12)



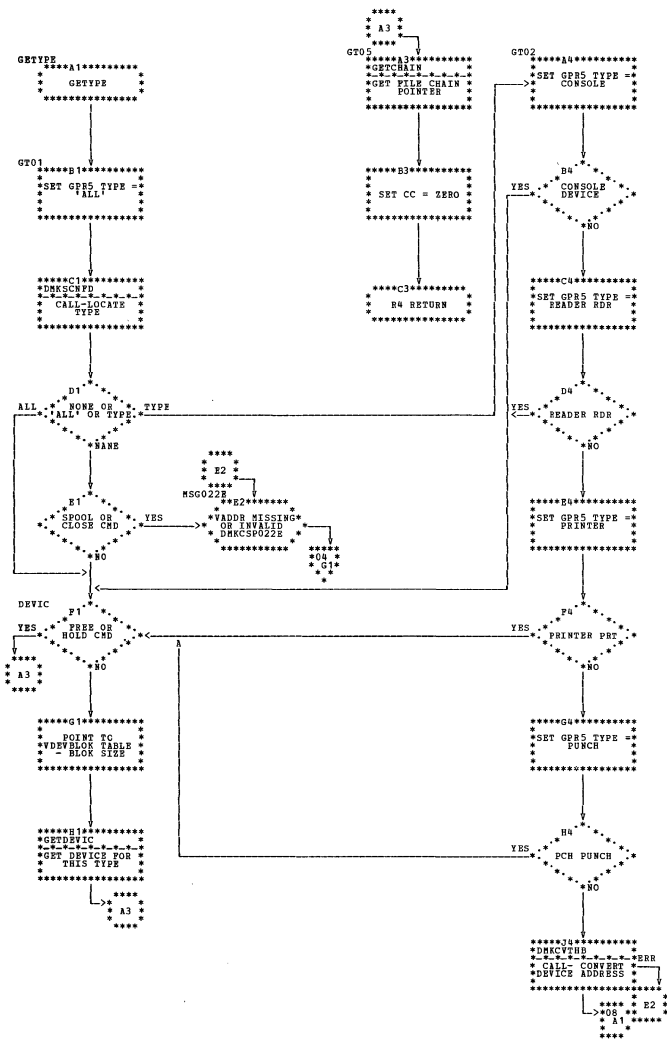


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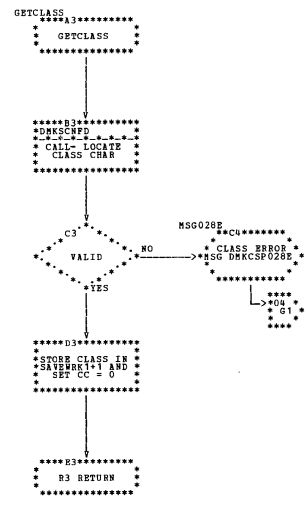
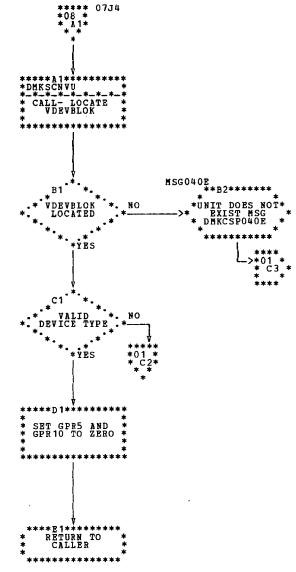


TO:G1  
1084  
1234



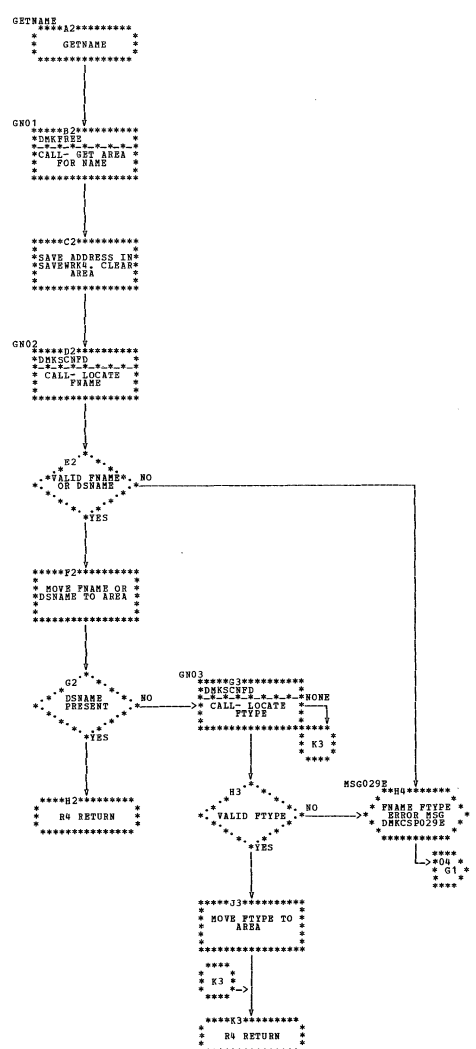
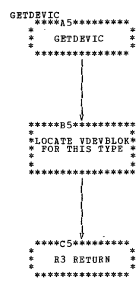
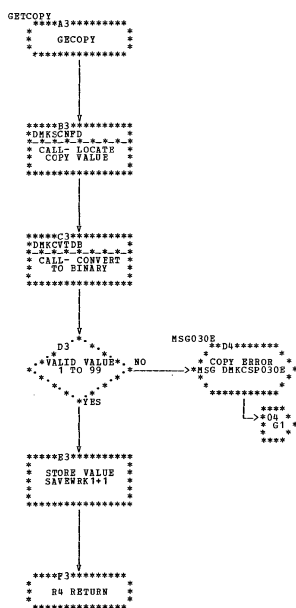
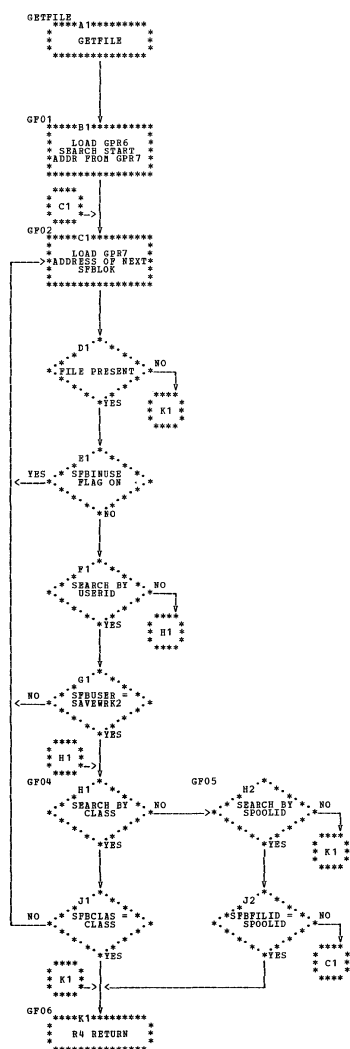


\* \* \* \* \*



| DMKCSPP -- Process Spooling Commands (Parts 7 and 8 of 12)

| DMKCSP -- Process Spooling Commands (Parts 9 and 10 of 12)



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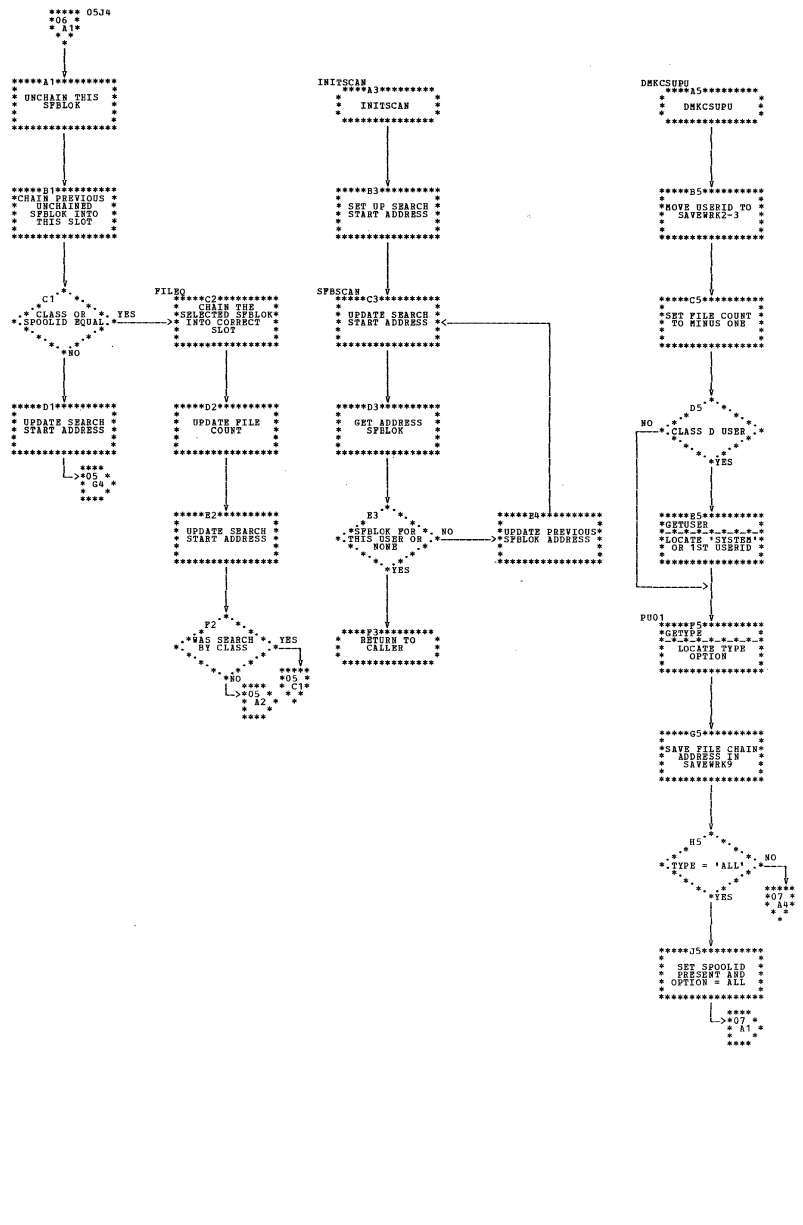
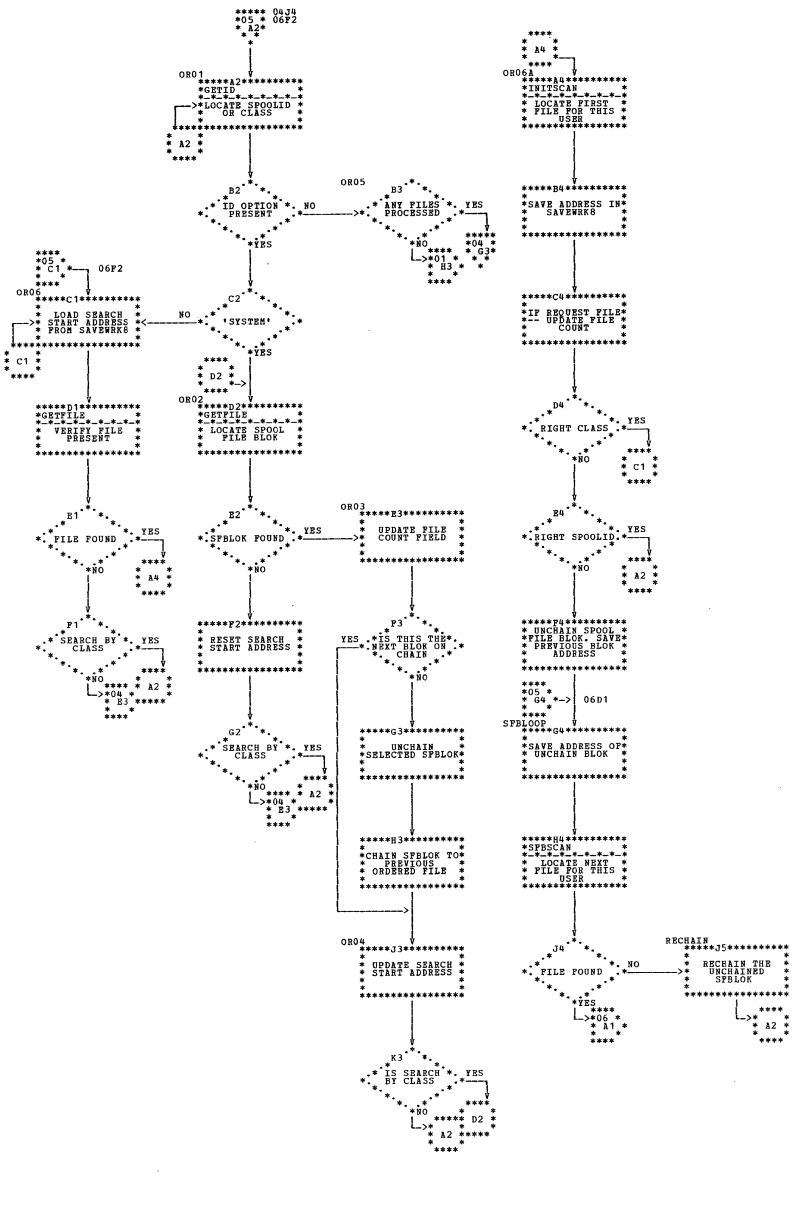








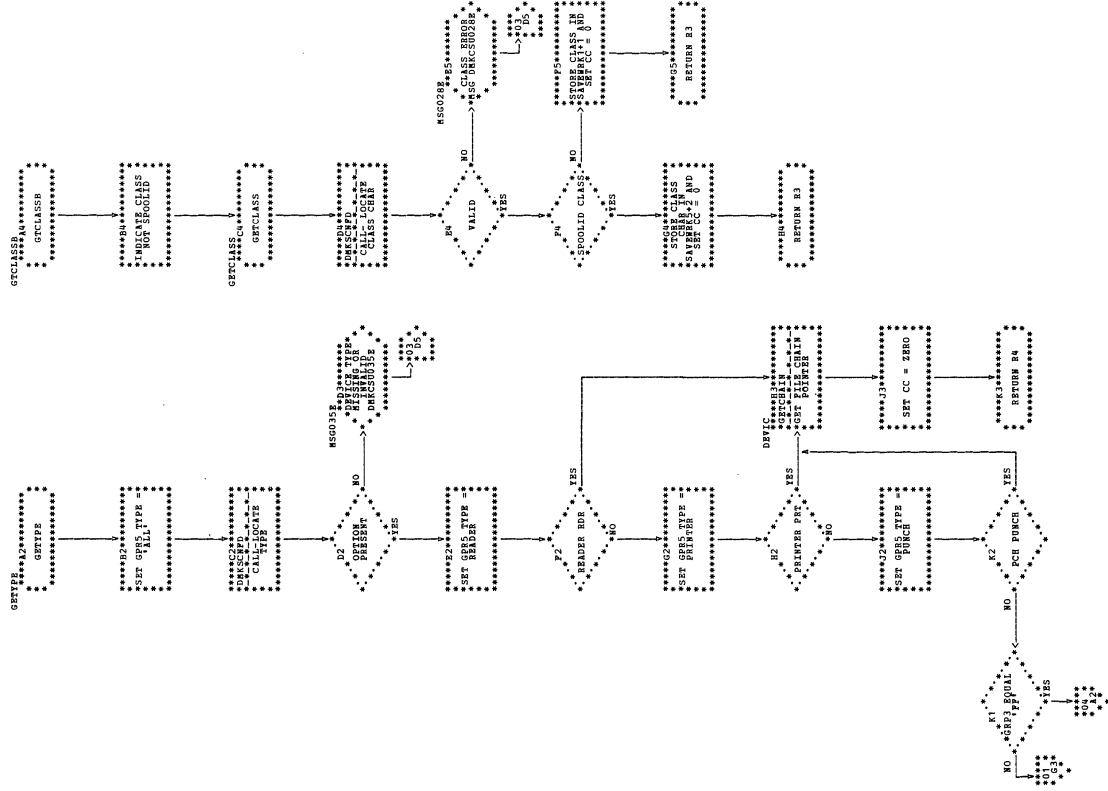
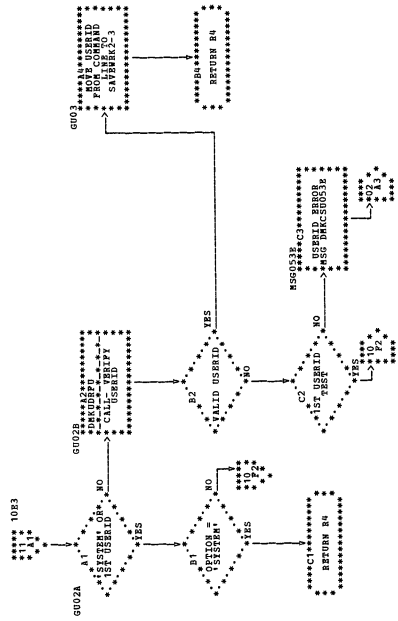
DMKCSU -- Process Spooling Commands (Parts 5 and 6 of 17)





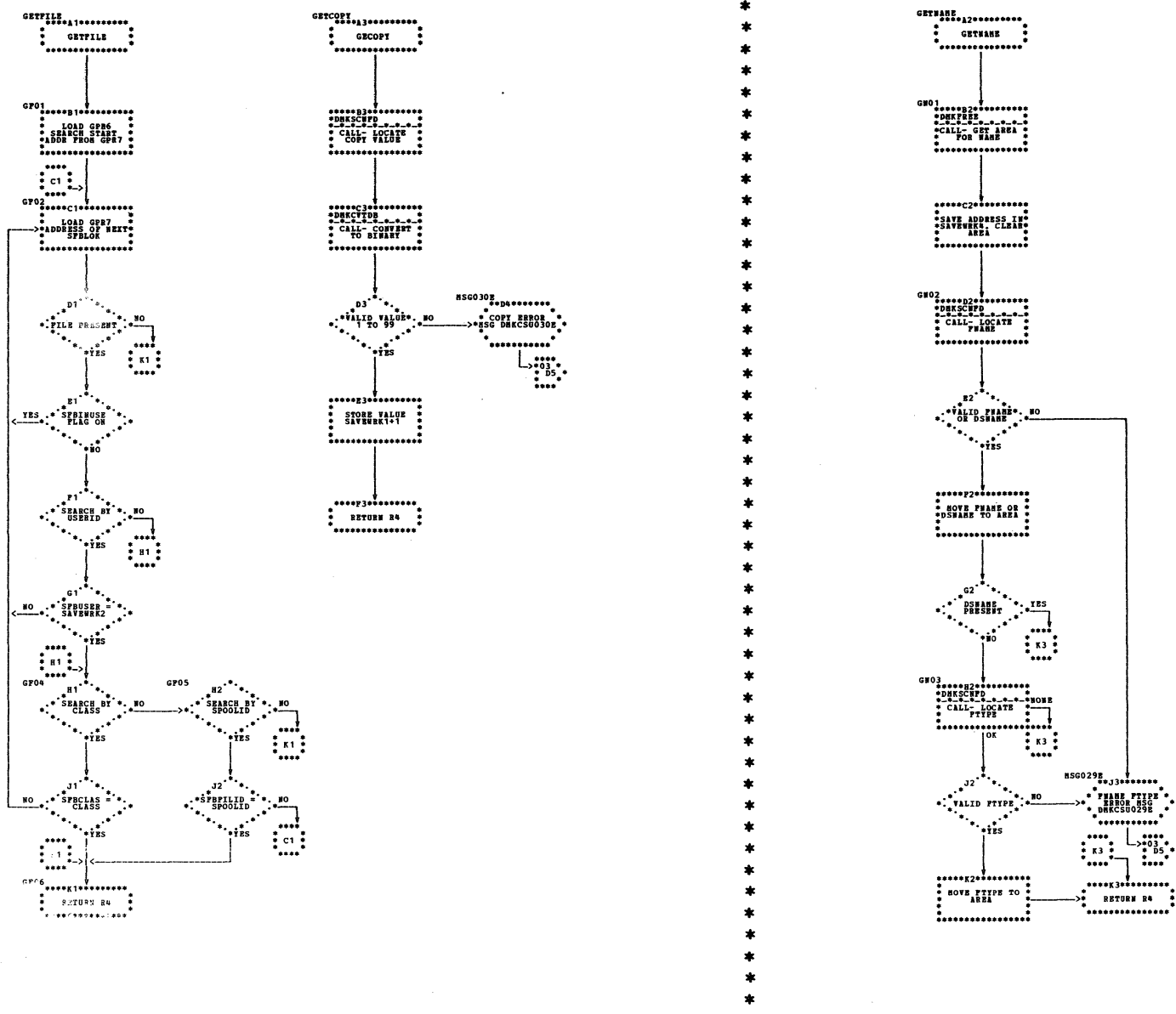


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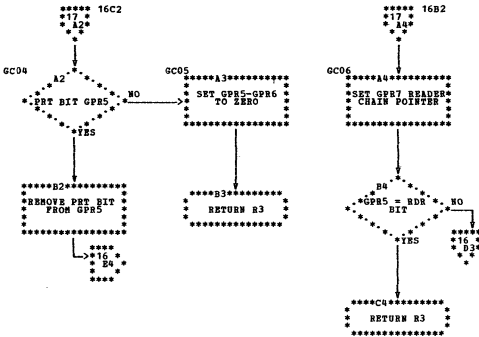
DMKCSU -- Process Spooling Commands (parts 11 and 12 of 17)

DMKCSU -- Process Spooling Commands (Parts 13 and 14 of 17)





DMKCSU -- Process Spooling Commands (Part 17 of 17)



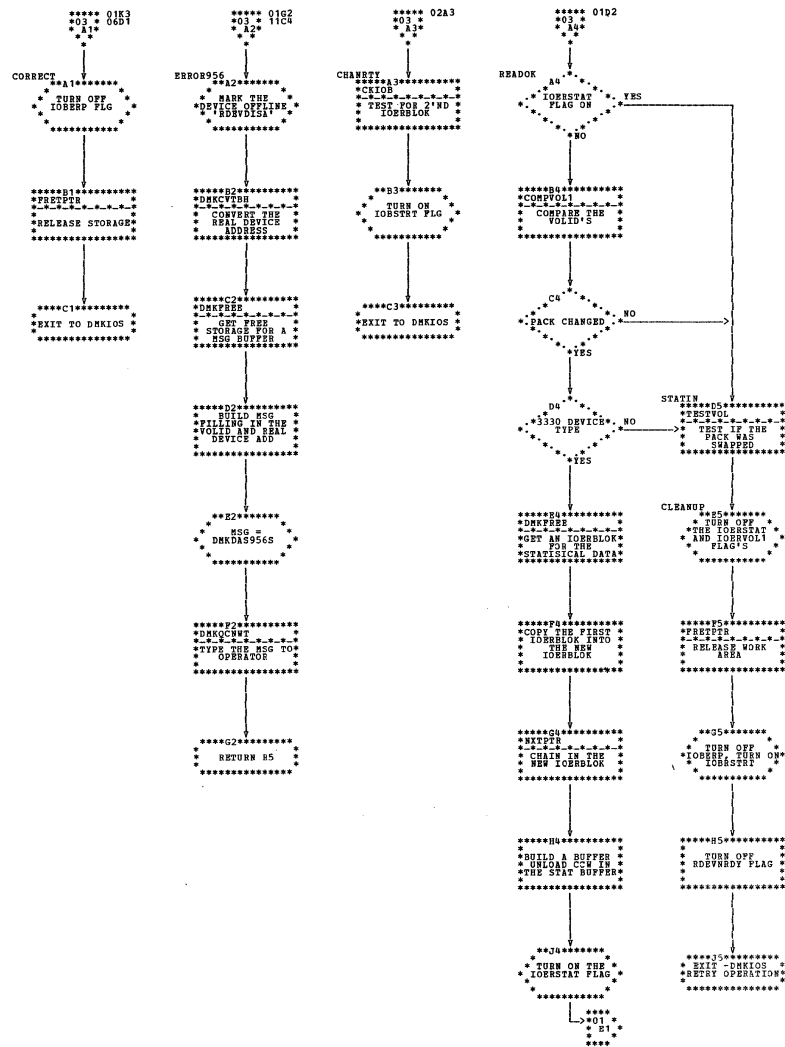




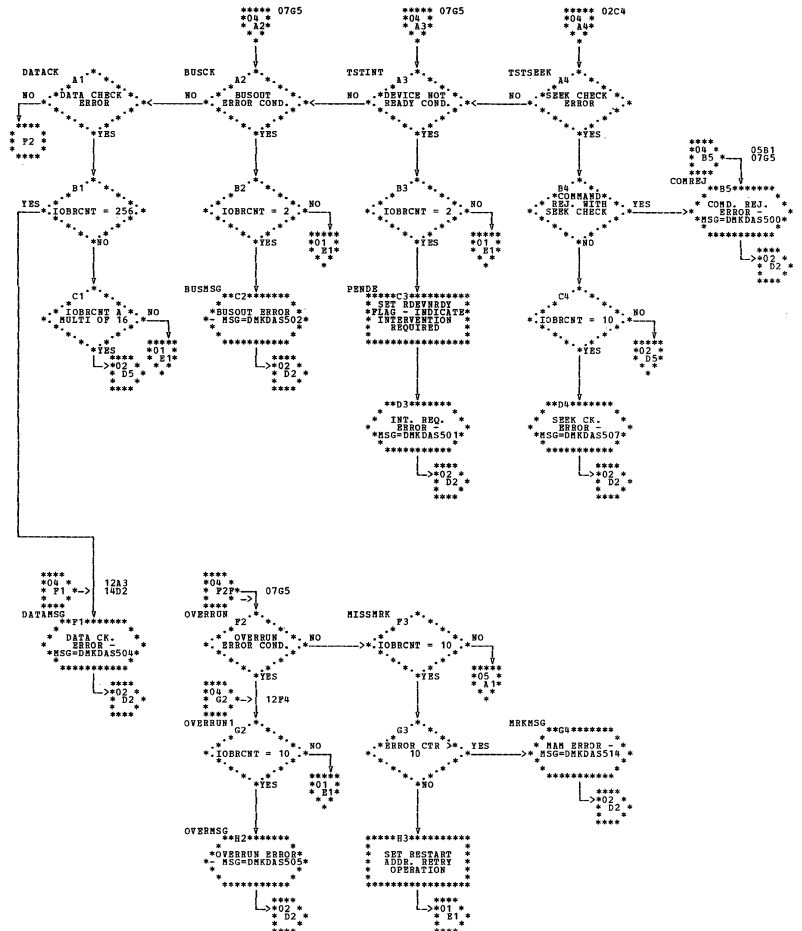




DMKDAS -- DASD Error Recovery Procedures (Parts 3 and 4 of 14)

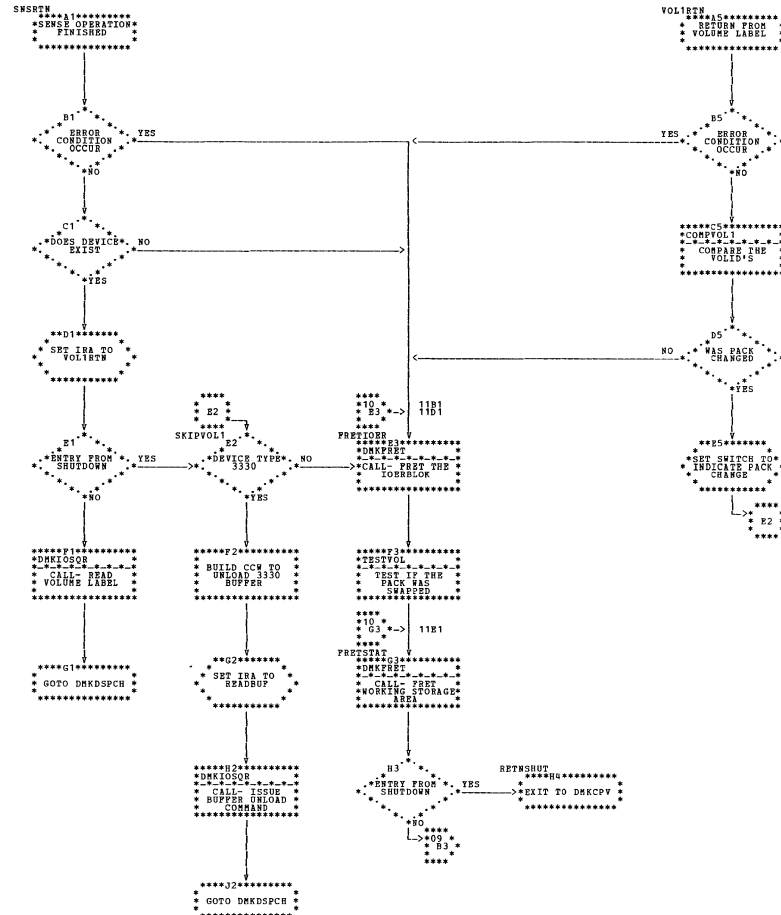
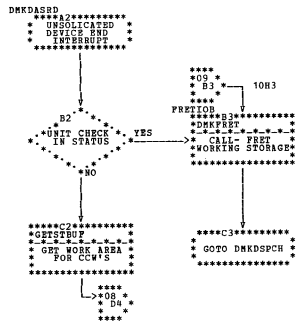


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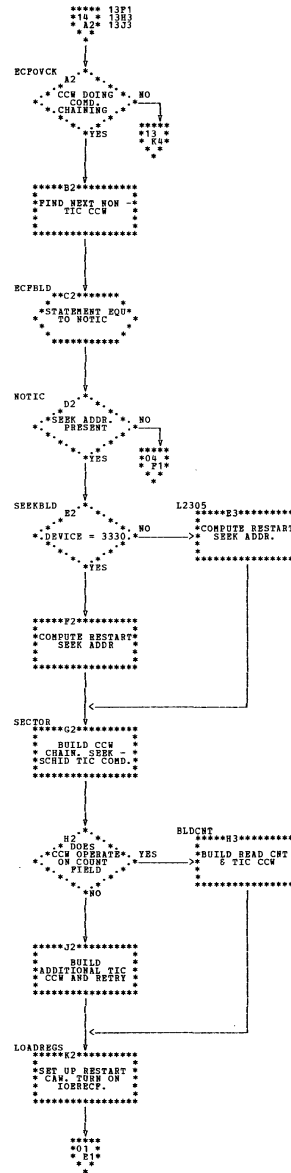
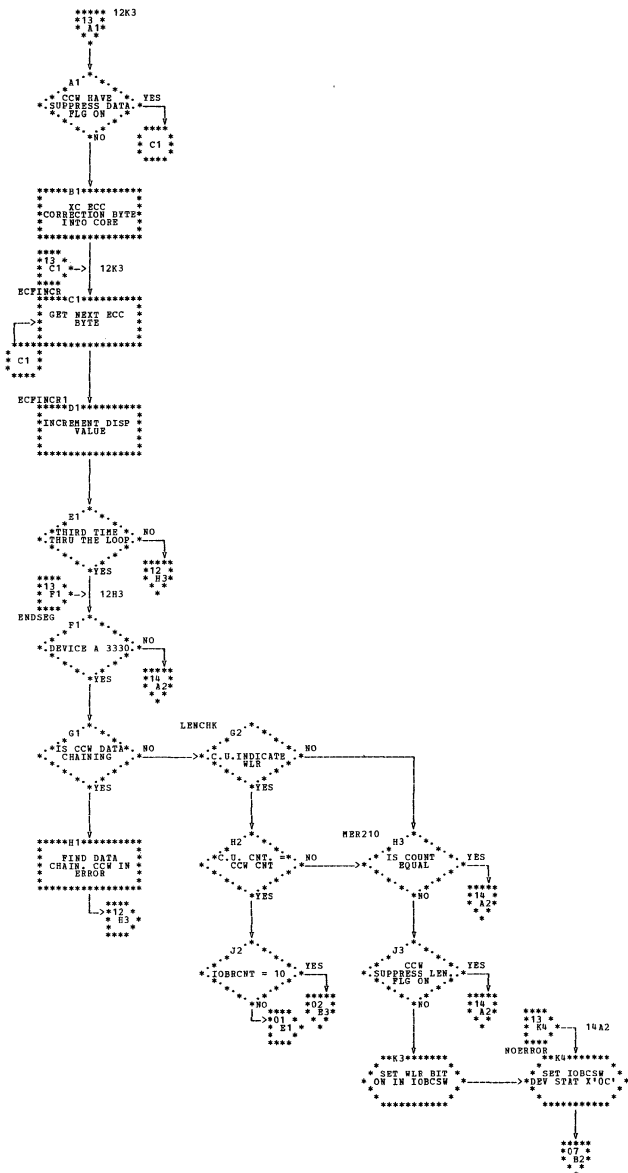




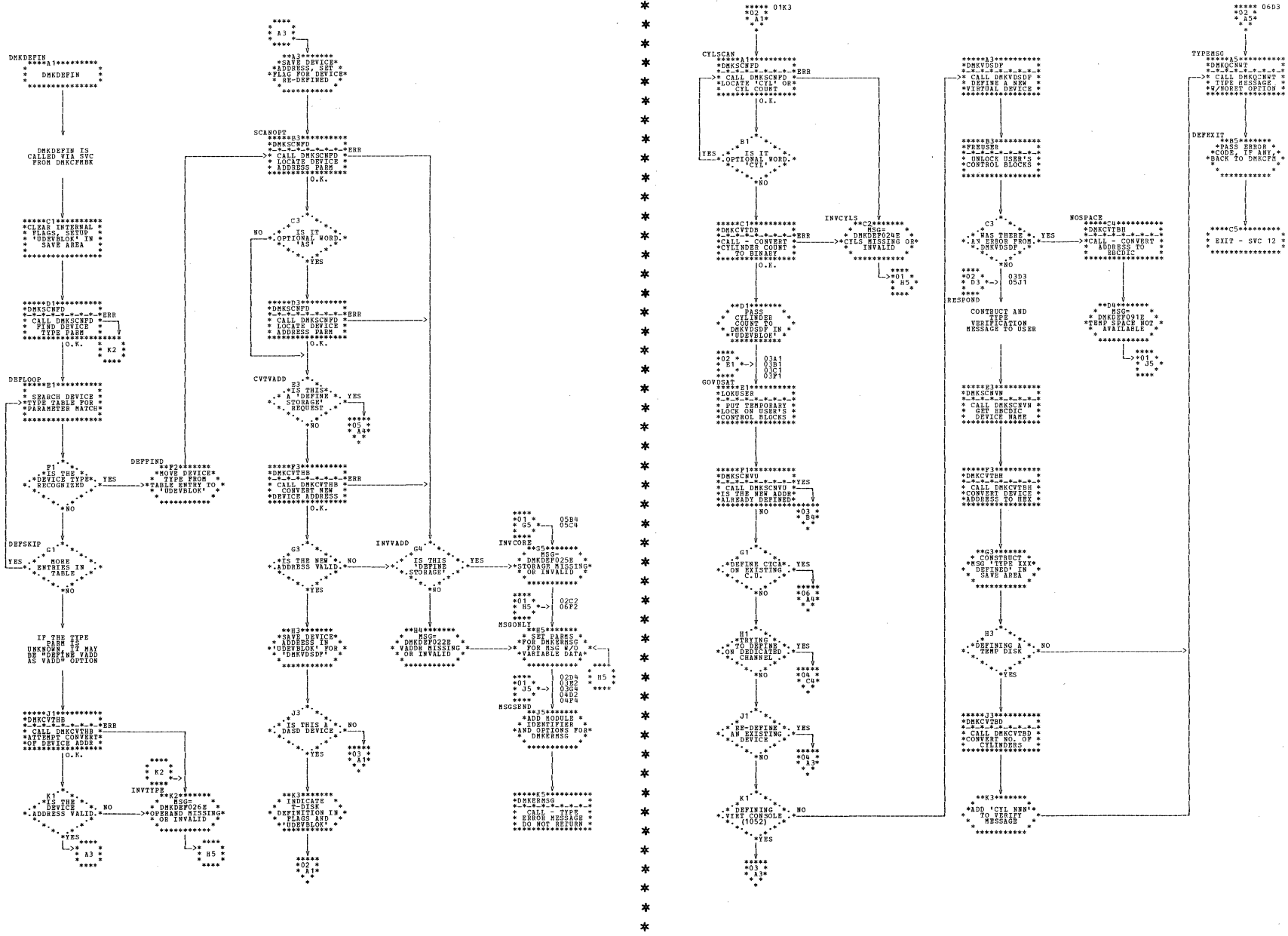






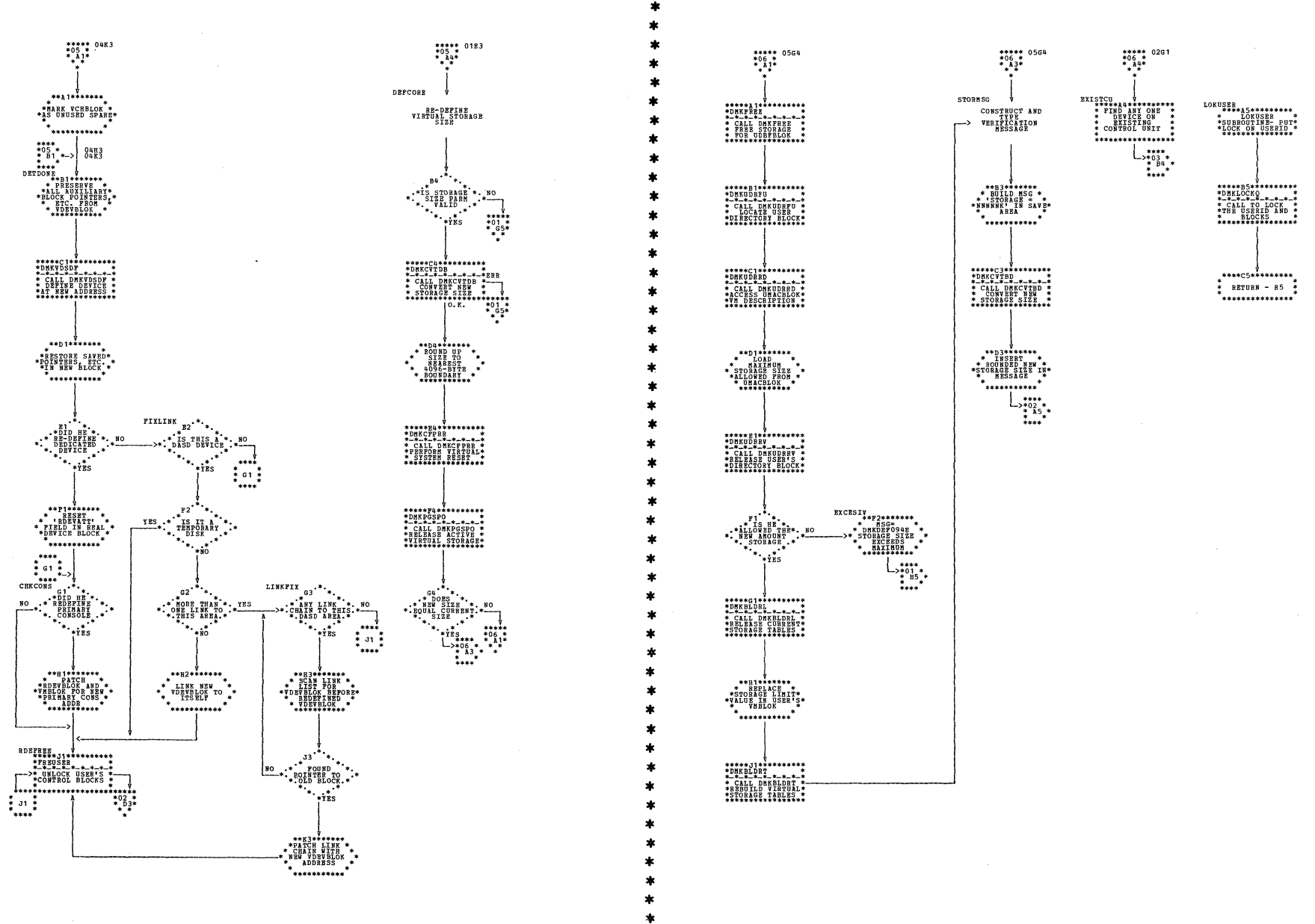


| DMKDFE -- Process DEFINE Command; Define Virtual Storage or a Device (Parts 1 and 2 of 7)



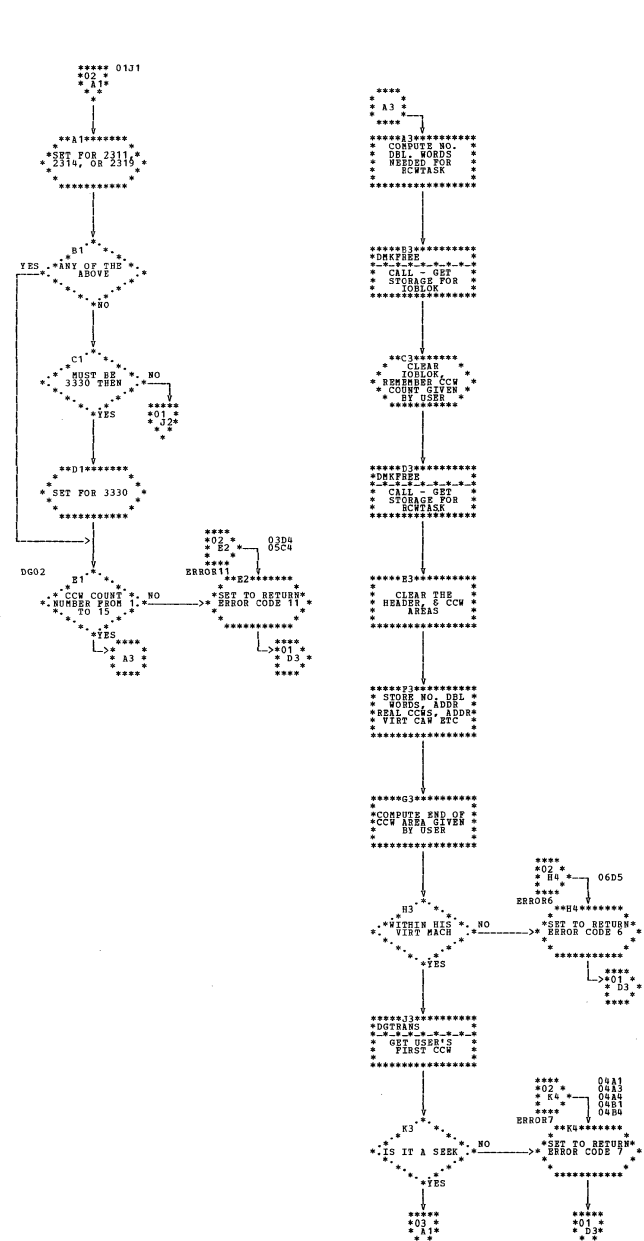
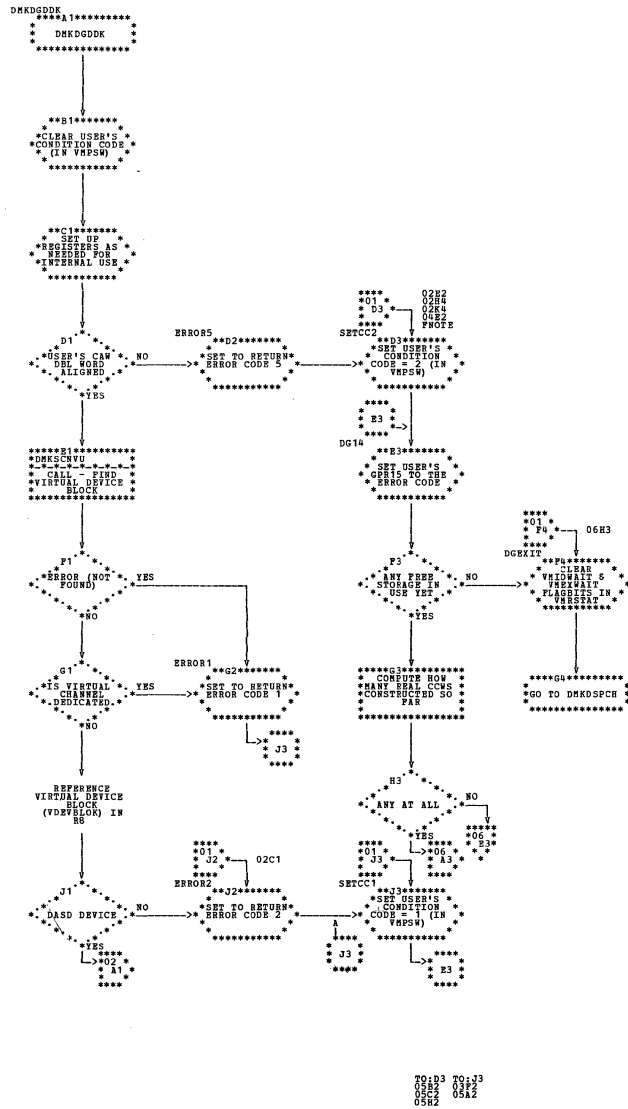


| DMKDEF -- Process DEFINE Command; Define Virtual Storage or a Device (Parts 5 and 6 of 7)

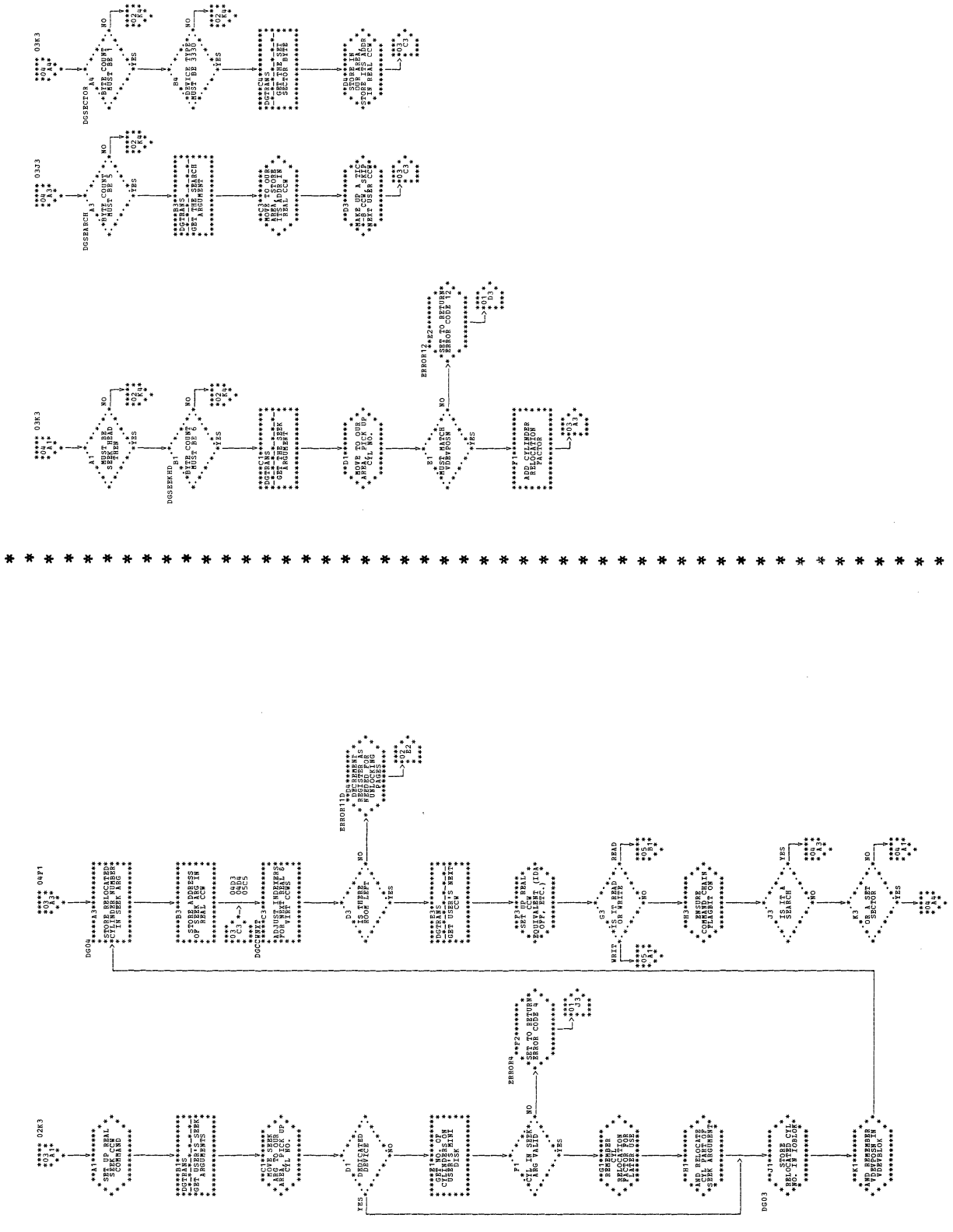




DMKDGDD -- Perform DASD I/O (Parts 1 and 2 of 8)



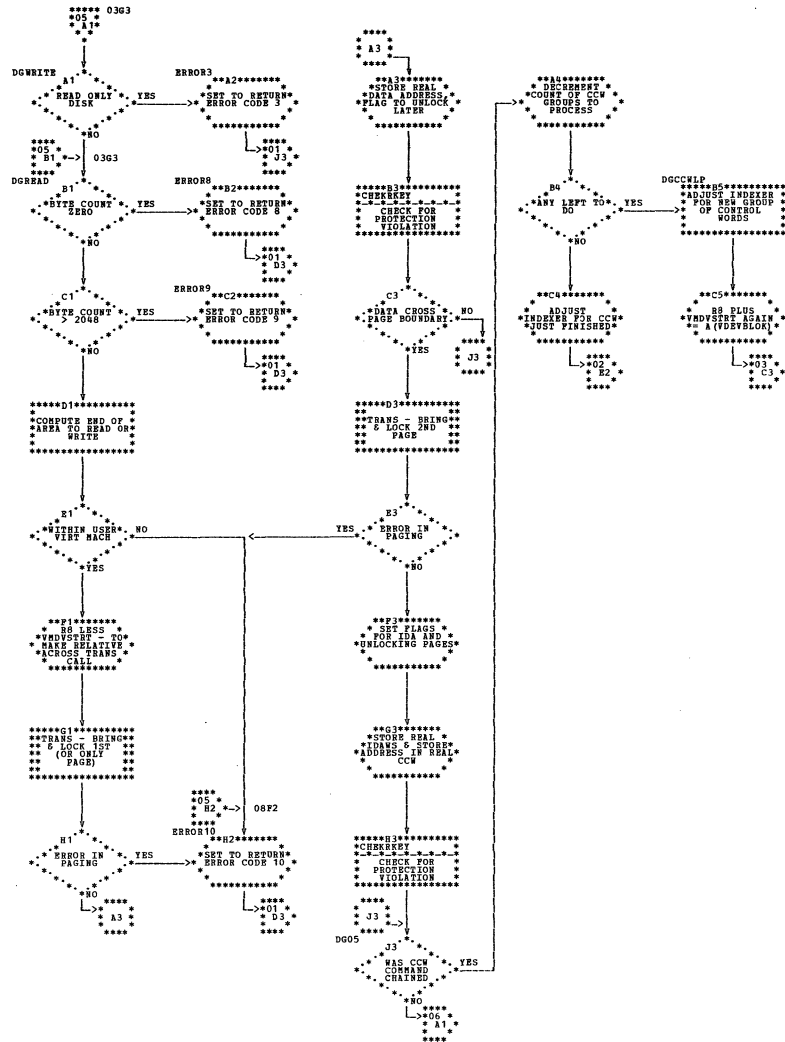
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 FROM: 03:29  
 CCB: 03A2  
 OSN2



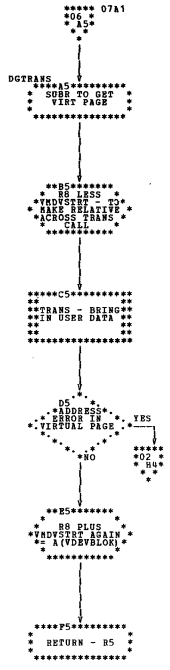
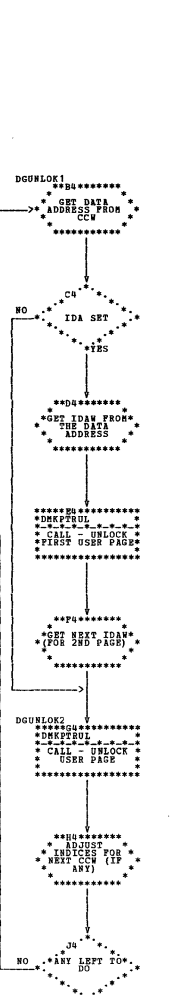
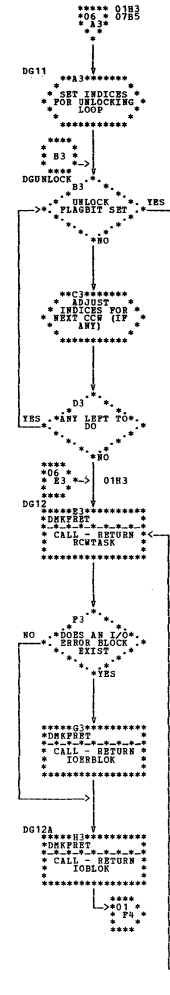
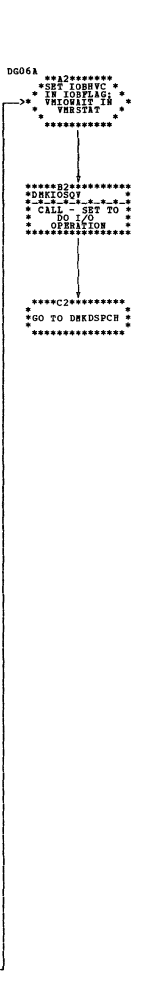
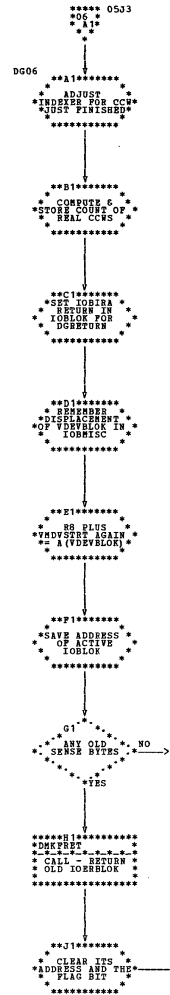
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I DMKDGD -- Perform DASD I/O (Parts 3 and 4 of 8)

| DMKGDG -- Perform DASD I/O (Parts 5 and 6 of 8)

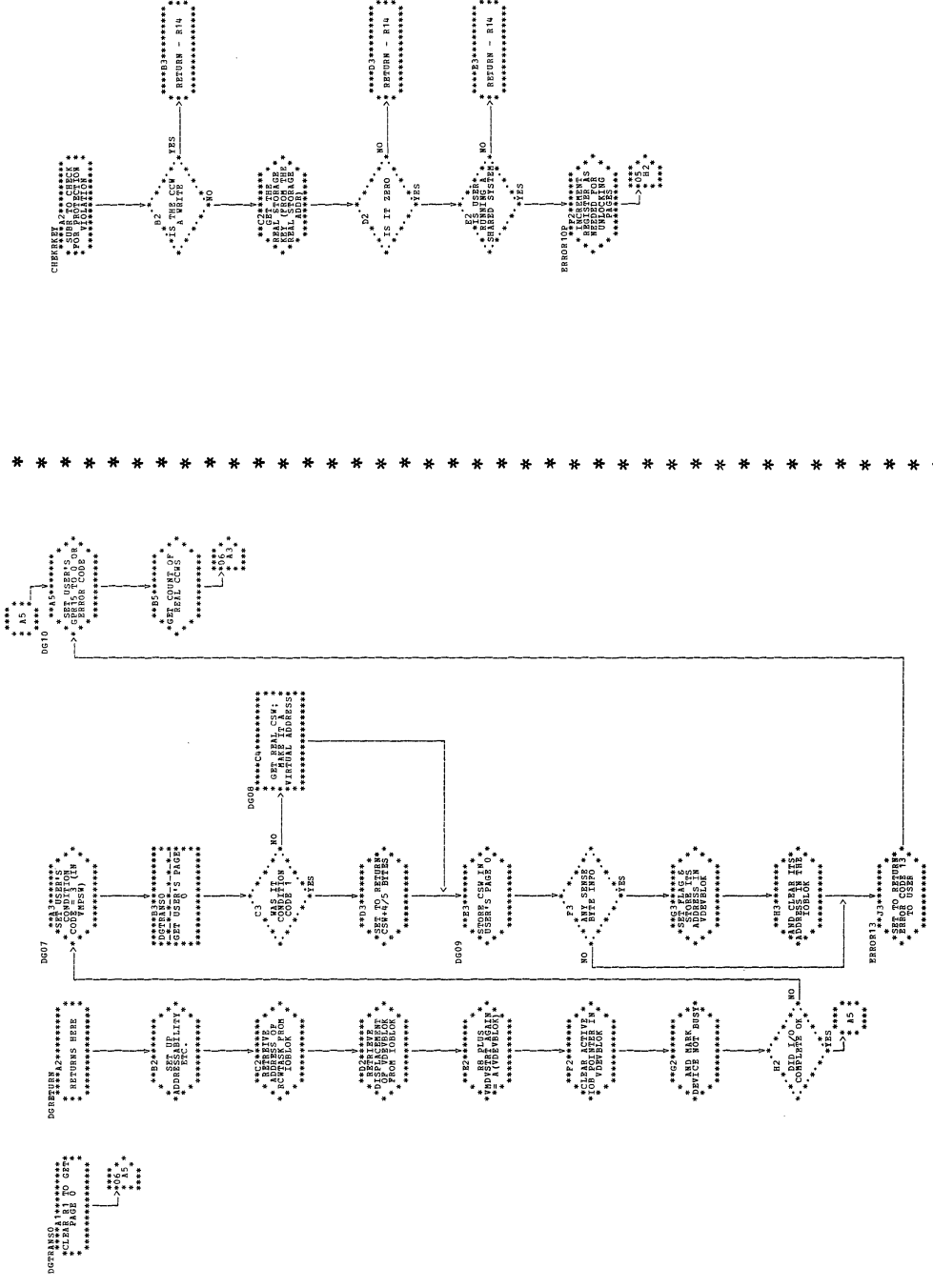


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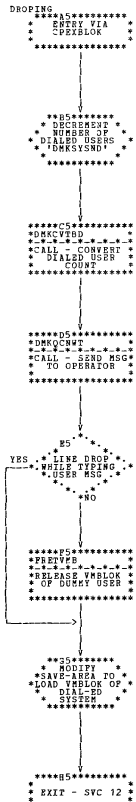
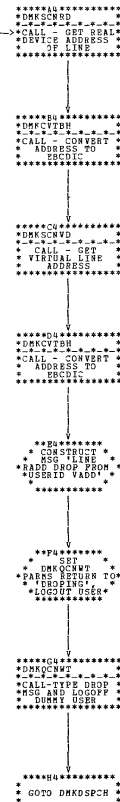
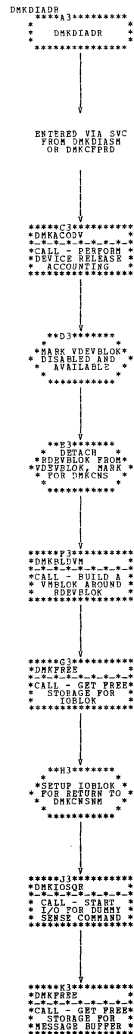
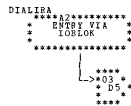
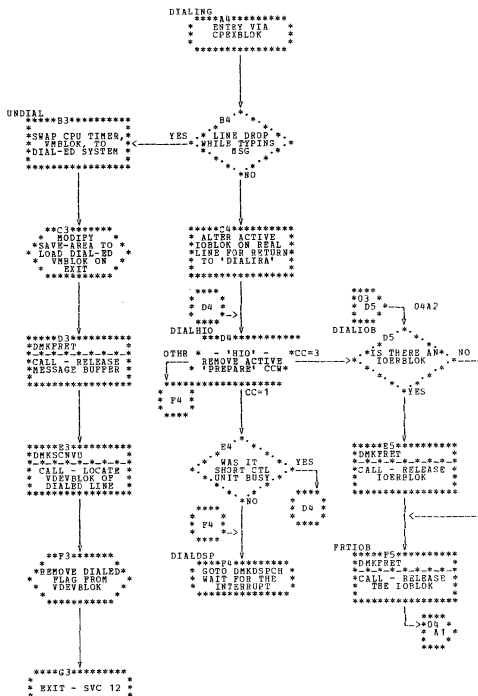
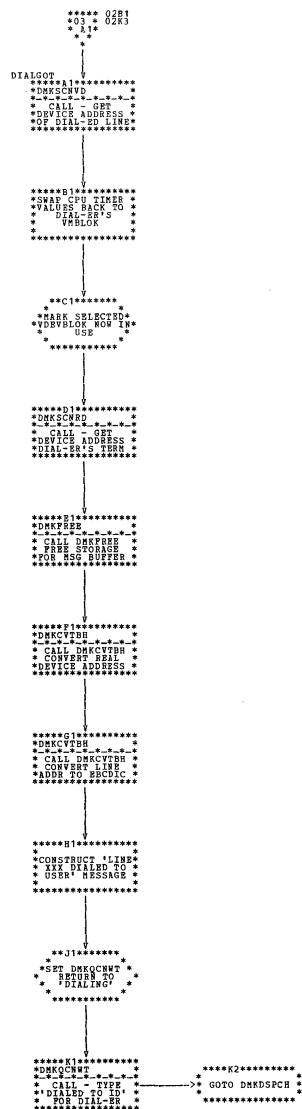


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I DMKGDG -- Perform DASD I/O (Parts 7 and 8 of 8)

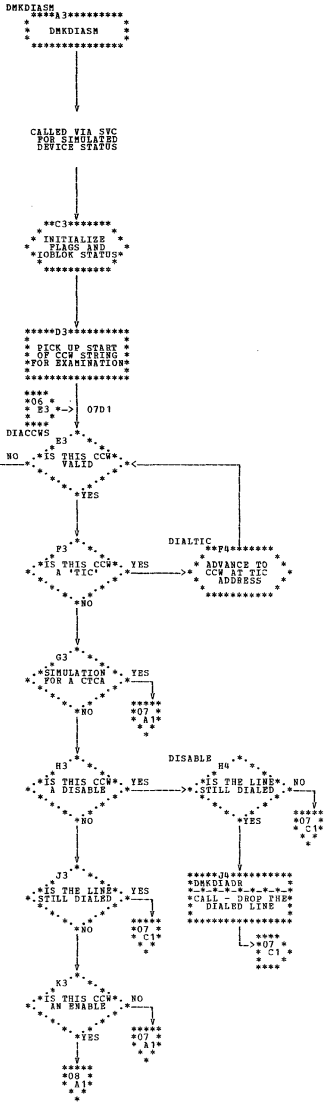
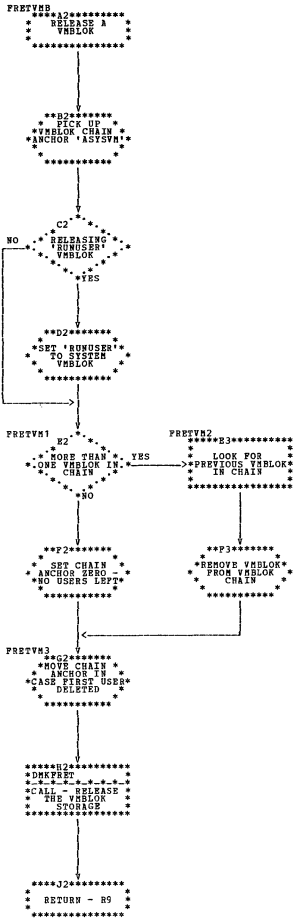


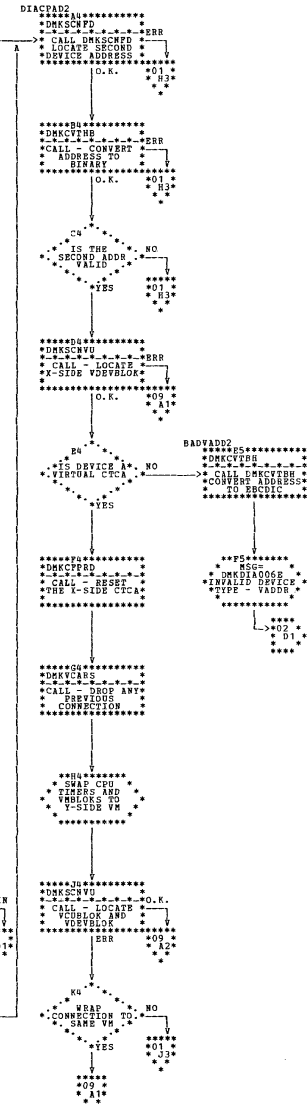
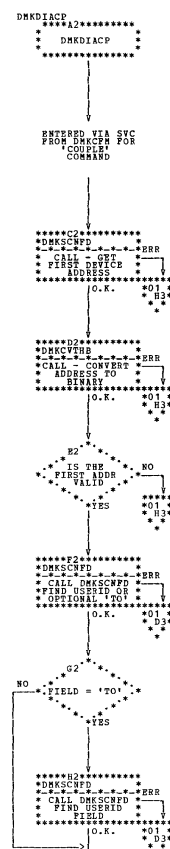
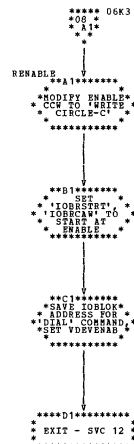
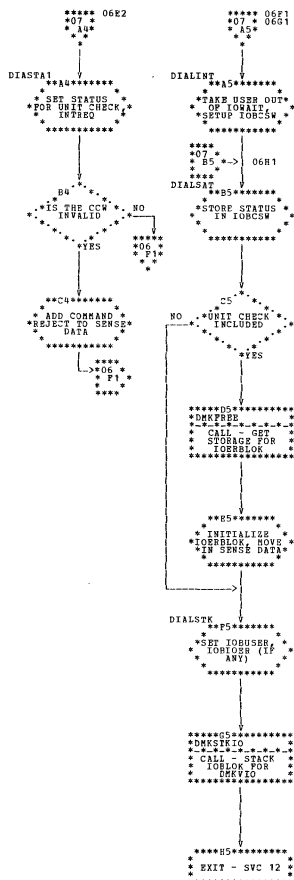
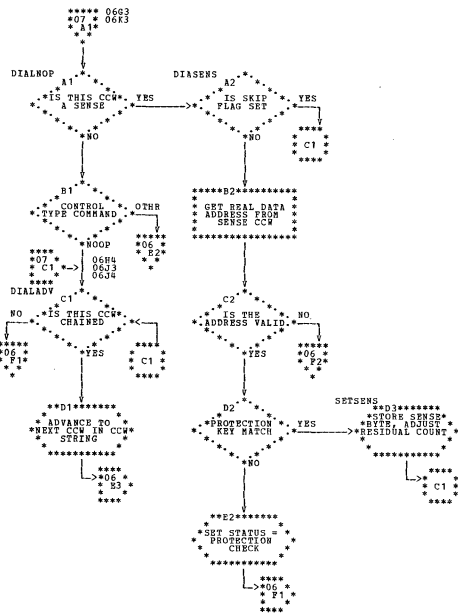


| DMKDIA -- Process DIAL Command, Connect to a Virtual System (Parts 3 and 4 of 9)

DMKDIA -- Process DIAL Command, Connect to a Virtual System (Parts 5 and 6 of 9)

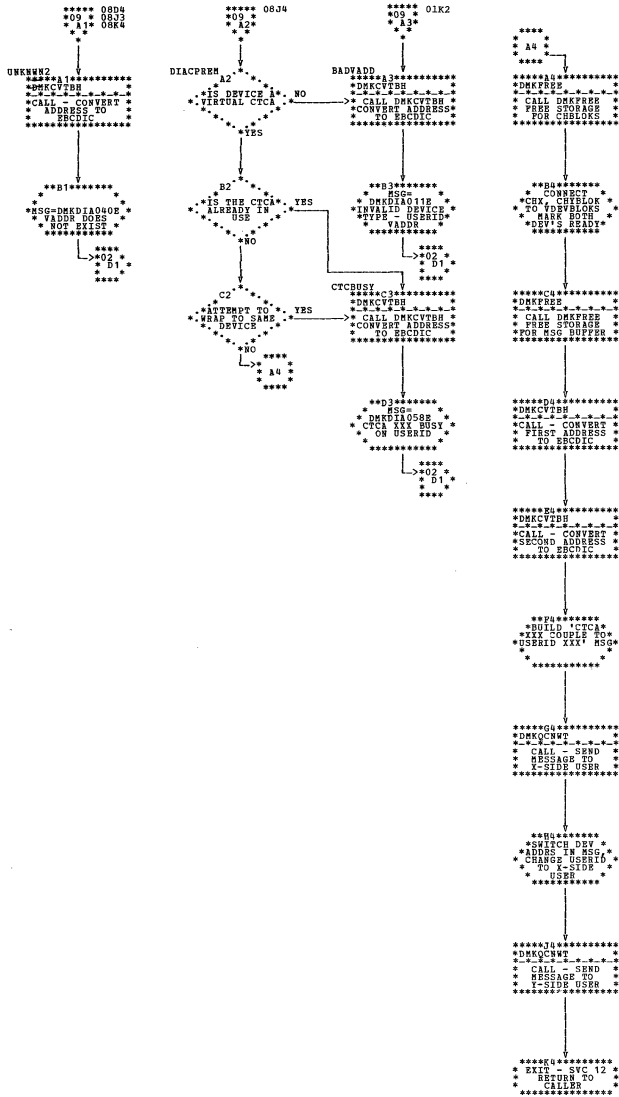
\* \* \* \* \*





| DMKDIA -- Process DIAL Command, Connect to a Virtual System (Parts 7 and 8 of 9)

| DMKDIA -- Process DIAL Command, Connect to a Virtual System (Part 9 of 9)



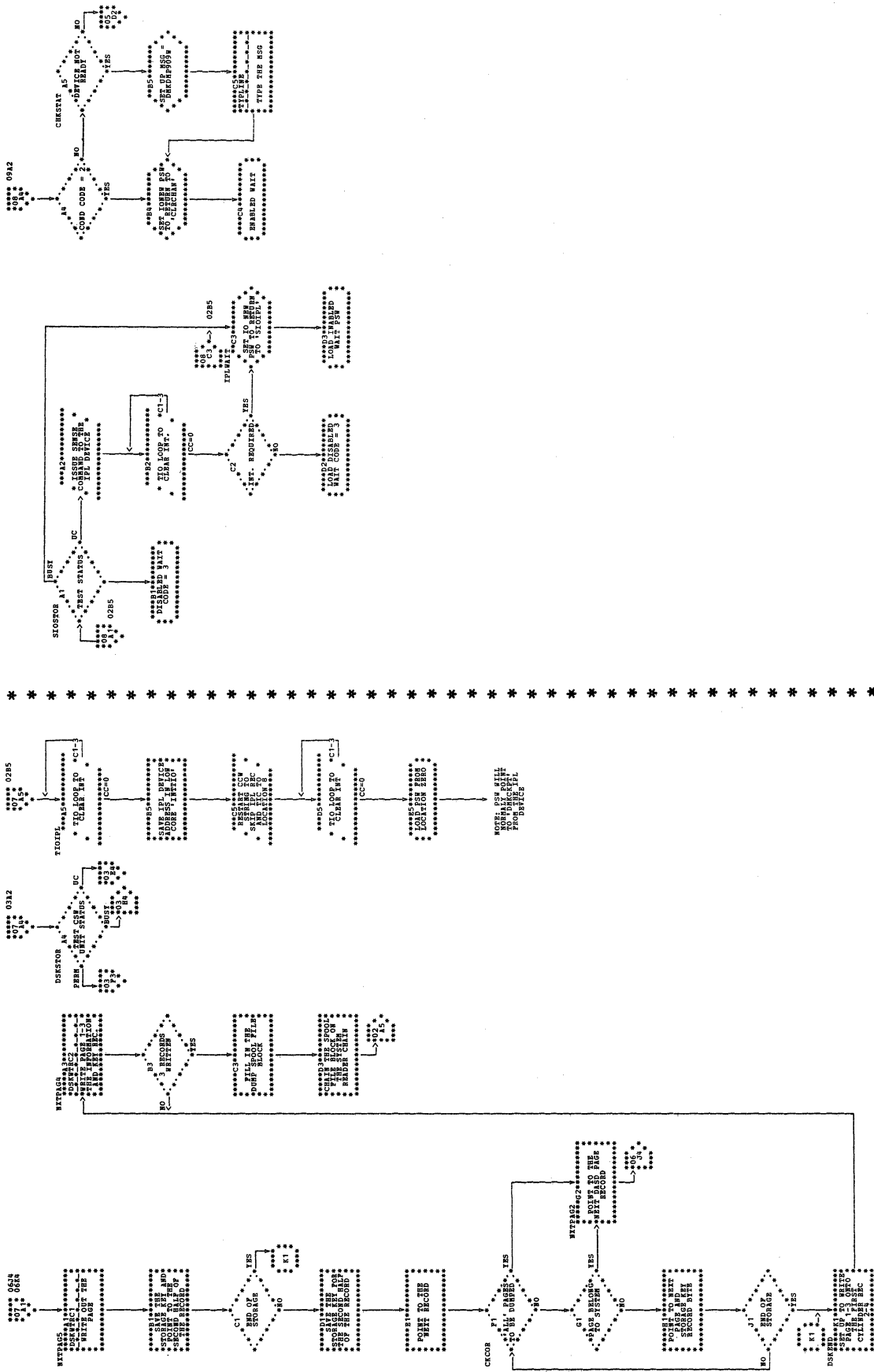




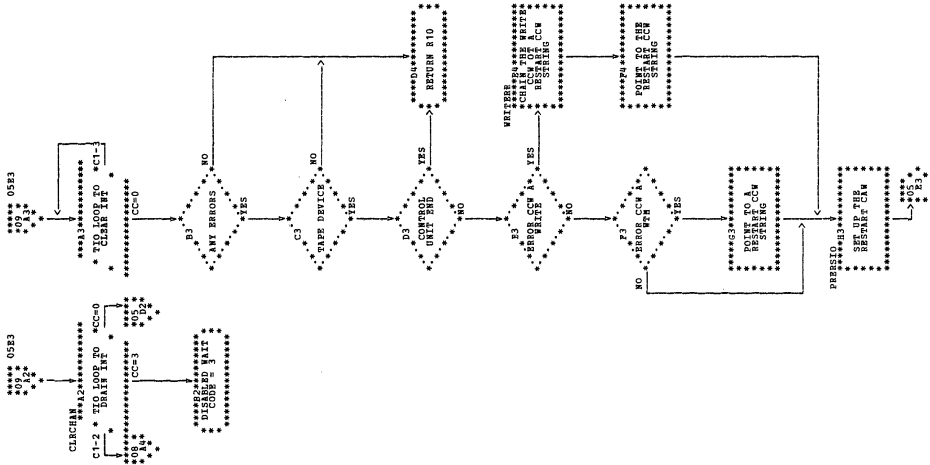




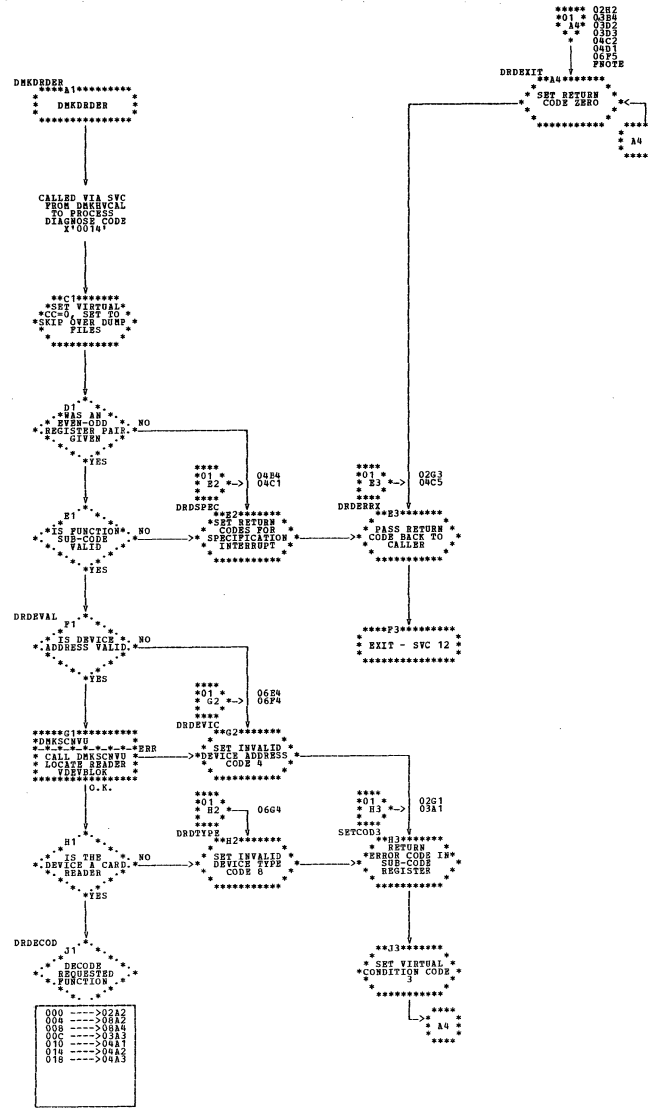
1 DDKDMP -- System Dump Processor (Parts 7 and 8 of 9)



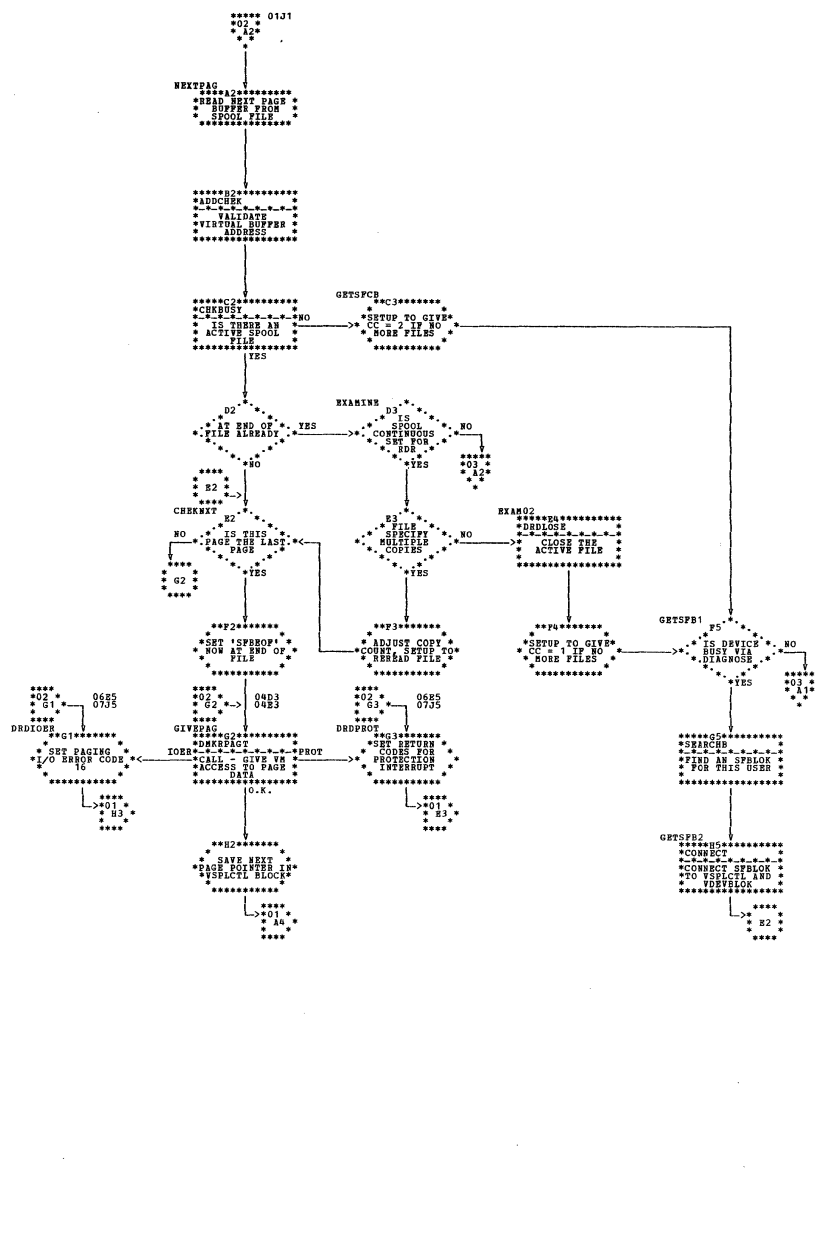
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DMKDRD -- Process Input Spool File (Parts 1 and 2 of 8)

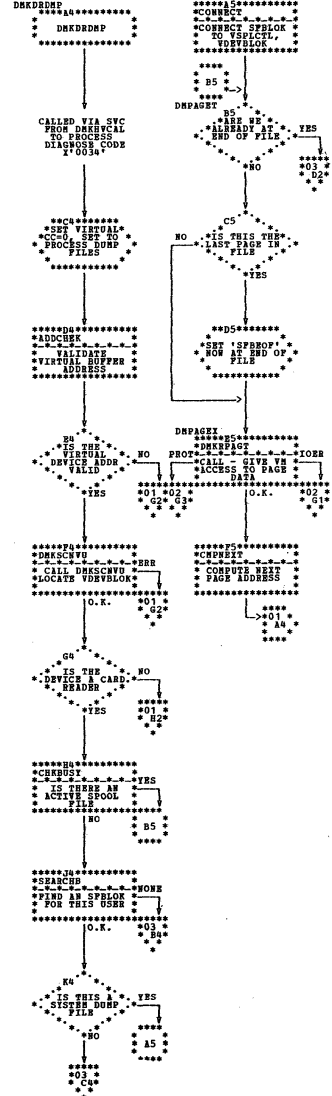
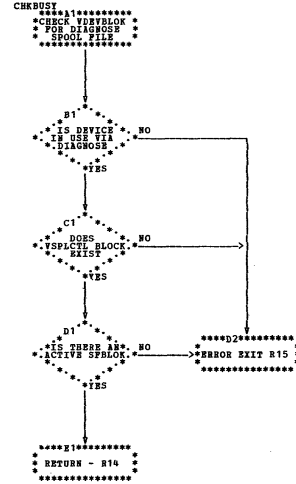
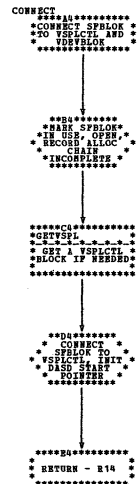
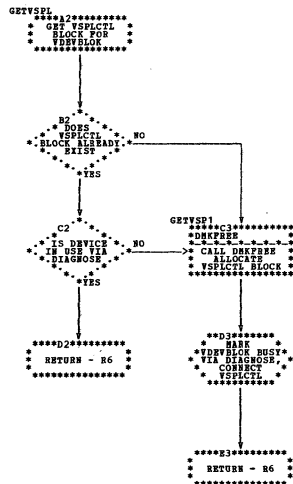
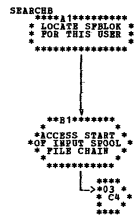


TO:14  
 0794  
 0795  
 08X3





DMKDRD -- Process Input Spool File (Parts 5 and 6 of 8)



\*\*\*\*\*

```

CHPNEXT *****
*CALCULATE DASD *
*ADDRESS OF NEXT *
*PAGE *
*****

```

```

*****B1*****
*ACCESS OWNED *
*VOLUME LIST TO *
*FIND RDEVBLK *
*****

```

```

*****C1*****
*FIND NUMBER OF *
*PAGES PER *
*CYLINDER FOR *
*THIS DEVICE *
*TYPE *
*****

```

```

FNDTYPE *****
**C2*****
*INCREMENT *
*PAGE NUMBER IN *
*DASD ADDRESS *
*****

```

```

D2 *****
*AT END OF *
*THIS CYLINDER *
*****

```

```

NEXTCYL *****
**D3*****
*MOVE TO PAGE *
*ONE OR TWO *
*CYLINDER *
*****

```

```

*****E2*****
RETURN - R7

```

```

*****E3*****
RETURN - R7

```

```

DMKDRDD *****
*DMKDRDD *
*****

```

```

*****C4*****
*DMKDRDD IS *
*CALLED VIA SVC *
*TO DELETE *
*SYSTEM DUMP *
*SPOOL FILE *
*****

```

```

*****C4*****
*CONSTRUCT DUMMY *
*SWAPTABLE ENTRY *
*FOR 'DMKPGTSD' *
*****

```

```

*****D4*****
*DMKPGTSD *
*CALL DMKPGTSD *
*DELETE FIRST *
*PAGE RECORD *
*****

```

```

*****D4*****
*AT END OF *
*FILE ALREADY *
*****

```

```

*****E4*****
*CHPNEXT *
*COMPUTE NEXT *
*DASD PAGE *
*ADDRESS *
*****

```

```

DMKDEL *****
*DMKDEL *
*CALL DMKDEL *
*DELETE ONE PAGE *
*RECORD *
*****

```

```

*****H4*****
*END OF FILE *
*REACHED *
*****

```

```

DMKDRDS *****
*DMKDRDS *
*****

```

```

*****C5*****
*CALLED VIA SVC *
*FROM DAREVCAL *
*TO PROCESS *
*DIAGNOSE CODE *
*10038 *
*****

```

```

*****C5*****
*SET VIRTUAL *
*CONDITION CODE *
*TO ZERO *
*****

```

```

*****D5*****
*ADDCHEK *
*VALIDATE *
*VIRTUAL BUFFER *
*ADDRESS *
*****

```

```

*****D5*****
*ACCESS SYSTEM *
*WHOLE SECTOR *
*TO FIND ENTRY *
*FOR 'DMKSTATE' *
*****

```

```

*****F5*****
*IS THE *
*SCREEN *
*MARKED *
*AVAILABLE *
*****

```

```

*****G5*****
*ACCESS SYSTEM *
*SWAPTABLE ENTRY *
*FOR 'DMKSTATE' *
*****

```

```

*****H5*****
*DOES SYMBOL *
*TABLE EXIST *
*****

```

```

*****I5*****
*DMKPAGE *
*O.K. *
*CALL - GIVE *
*ACCESS TO *
*SYMBOL TABLE *
*****

```

\*\*\*\*\*

```

***** 01J1 *****
*00 *
*12 *
*****

```

```

PRTSFC *****
*SET INDICATOR *
*FOR PRINT SPOOL *
*FILE BLOCK *
*****

```

```

READSFE *****
**B2*****
*ADDCHEK *
*VALIDATE *
*VIRTUAL BUFFER *
*ADDRESS *
*****

```

```

*****C2*****
*SETUP TO GIVE *
*CC = IF *
*FILES FOUND *
*****

```

```

D2 *****
*IS THE *
*VIRTUAL *
*DEVICE BUSY *
*****

```

```

*****E2*****
*DEVICE *
*BUSY DUE TO *
*DIAGNOSE USE *
*****

```

```

*****F2*****
*CAN WE *
*PREPARE CHAIN *
*SEARCH *
*****

```

```

*****G2*****
*SETUP TO GIVE *
*CC = IF *
*MORE FILES *
*****

```

```

READSCH *****
**B3*****
*IS THIS *
*SPOOL THE *
*CORRECT TYPE *
*****

```

```

*****H3*****
*GETVSP *
*INITIALIZE *
*DEVICE FOR *
*DIAGNOSE READS *
*****

```

```

*****J3*****
*TRANS *
*RING IN *
*USER'S SFBLOCK *
*BUFFER AREA *
*****

```

```

*****K3*****
*MOVE SFBLOCK *
*TO VIRTUAL *
*STORAGE *
*****

```

```

***** 01J1 *****
*00 *
*14 *
*****

```

```

PCHSFC *****
*SET INDICATOR *
*FOR PUNCH SPOOL *
*FILE BLOCK *
*****

```

DMKDRD -- Process Input Spool File (Parts 7 and 8 of 8)

SY20-0880-1, Page Modified by TNL SN20-2624, August 15, 1973

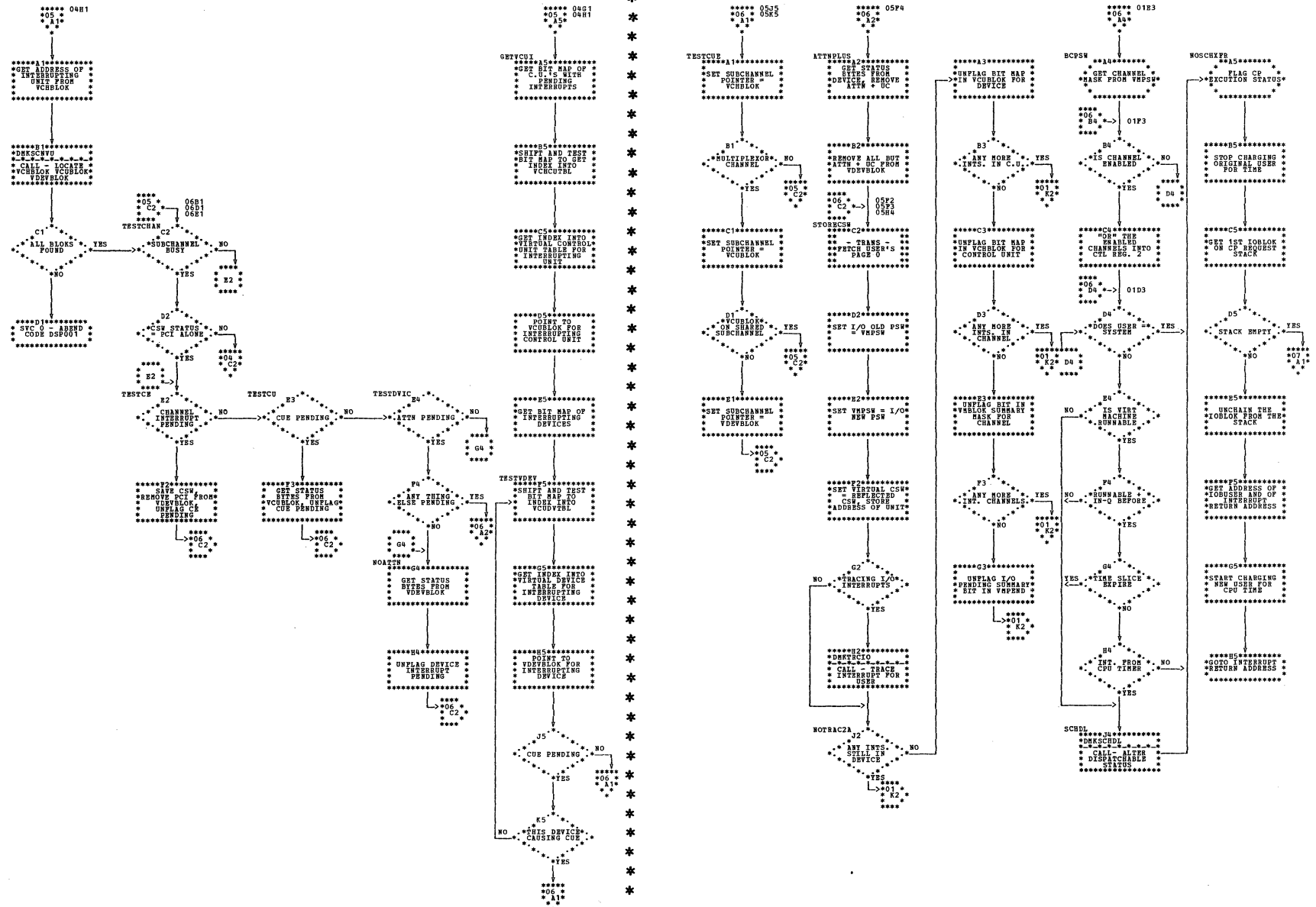
Program Organization 259

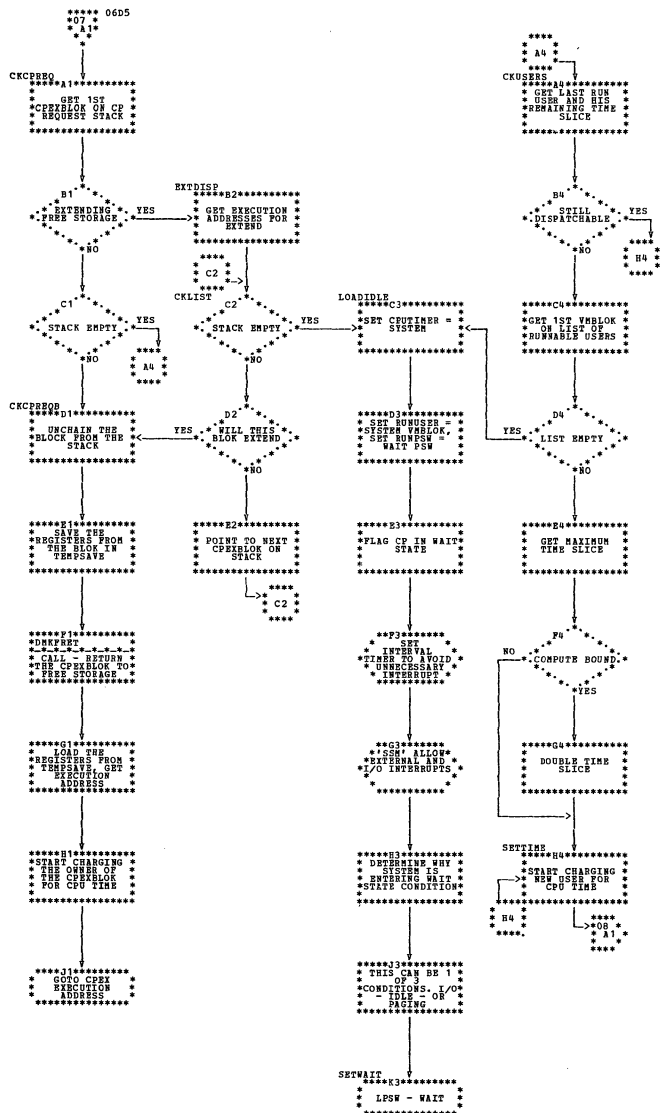




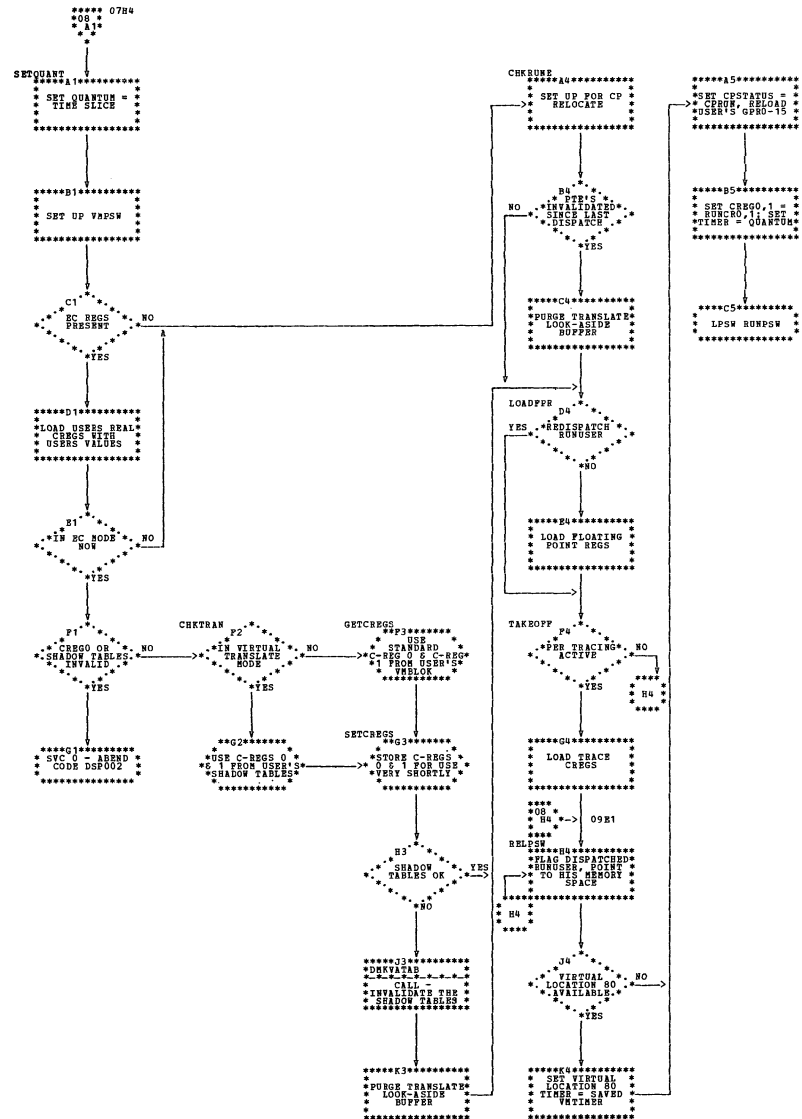


DMKDSP -- Dispatcher (Parts 5 and 6 of 10)



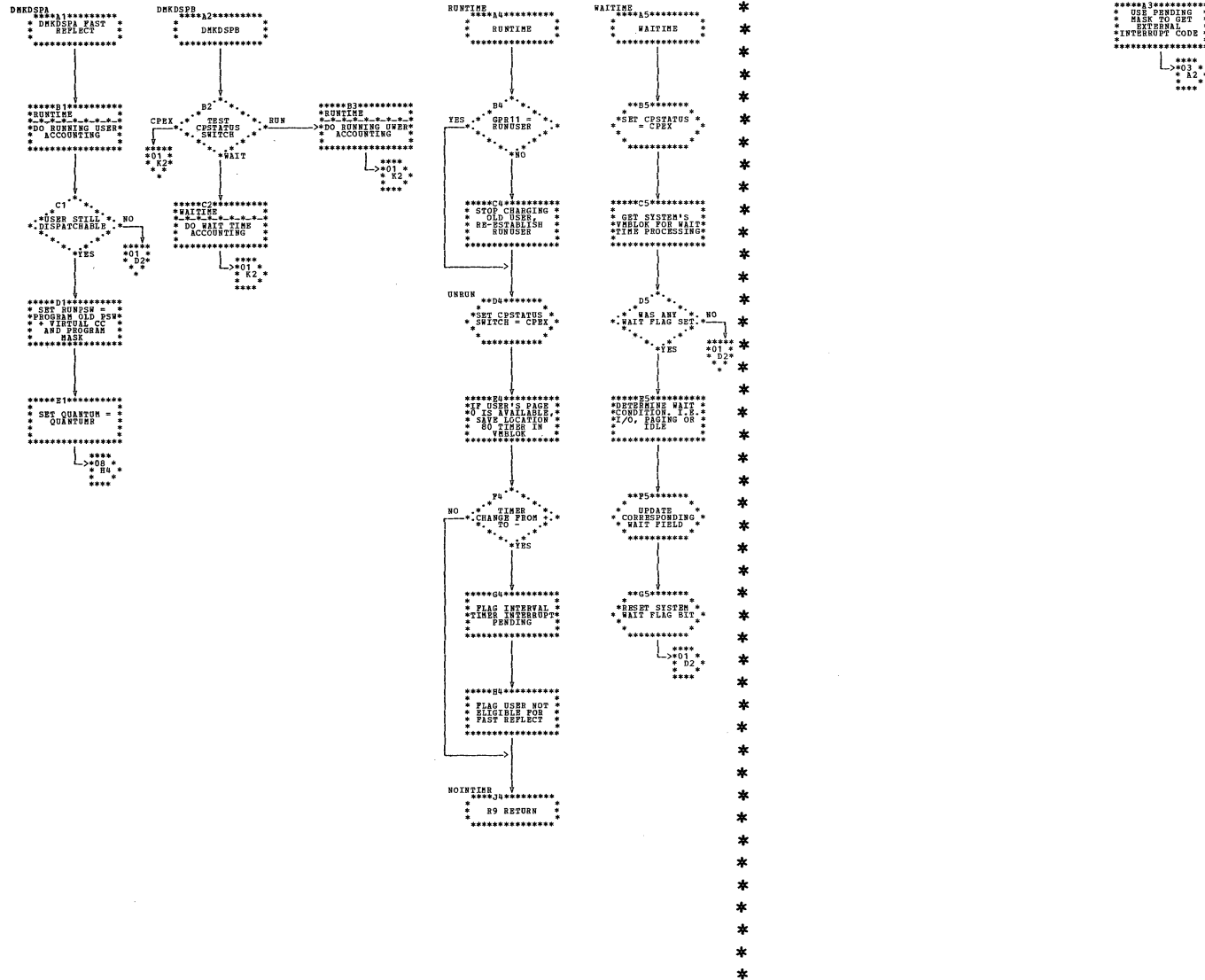


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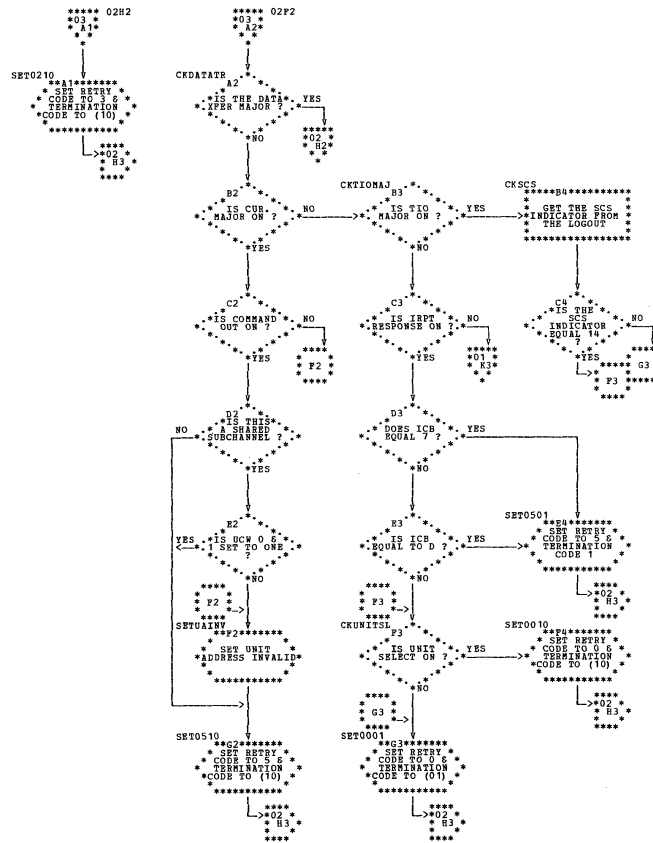
| DMKDSP -- Dispatcher (Parts 7 and 8 of 10)

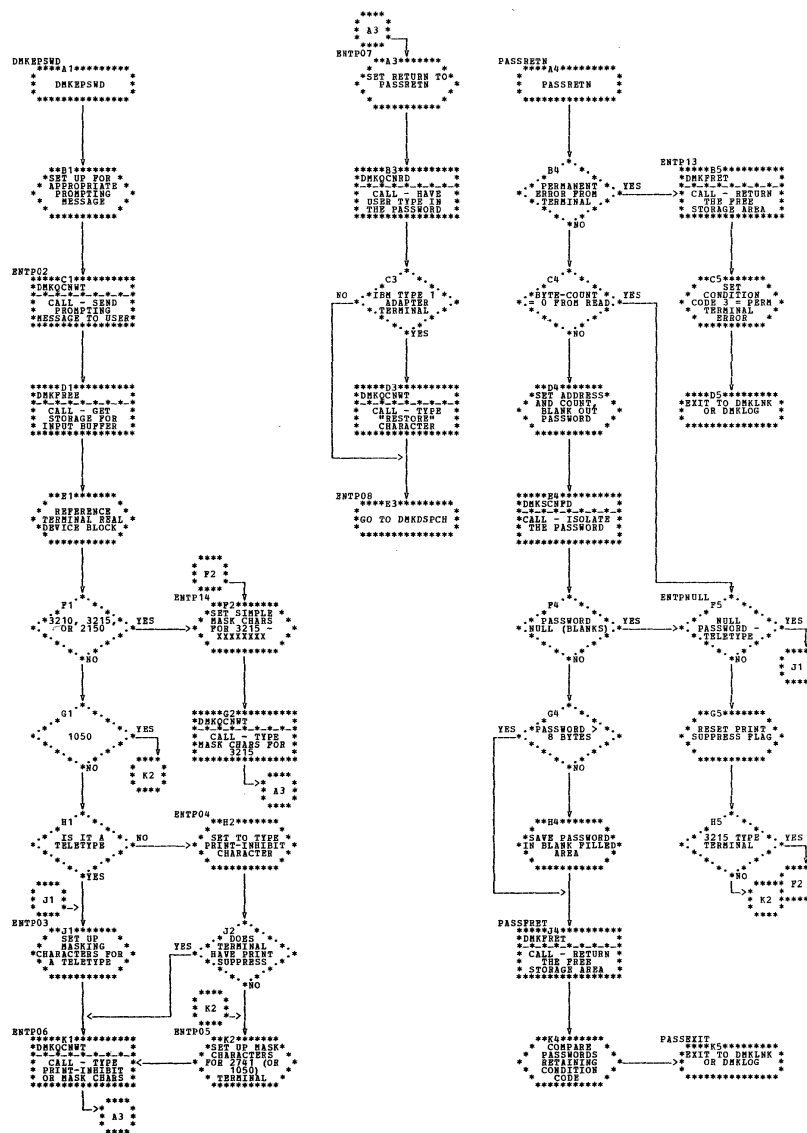
| DMKDSP -- Dispatcher (Parts 9 and 10 of 10)



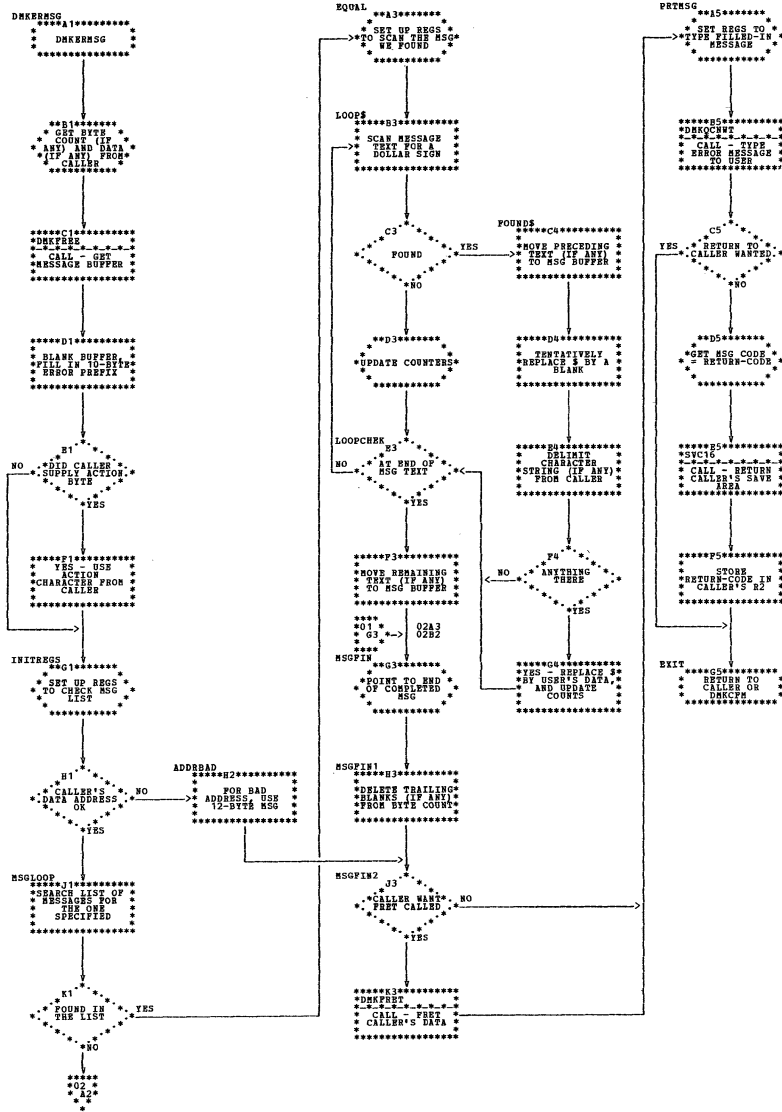


| DMKEIG -- 2880 Channel Module (Part 3 of 3)

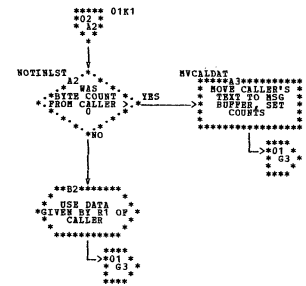




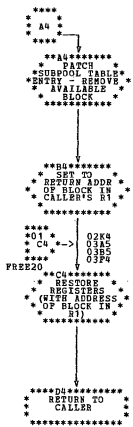
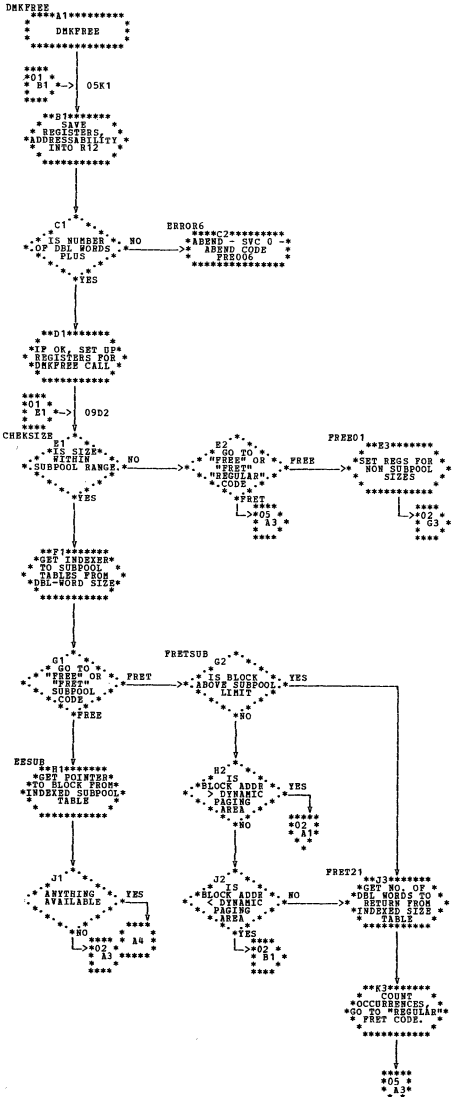
DMKERM -- Message Writer (Parts 1 and 2 of 2)



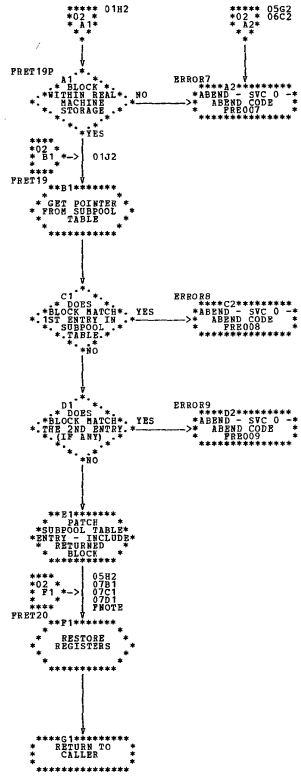
\* \* \* \* \*



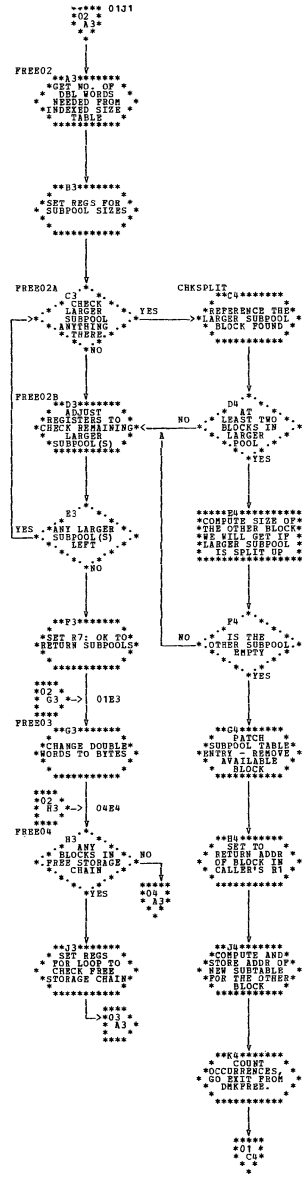




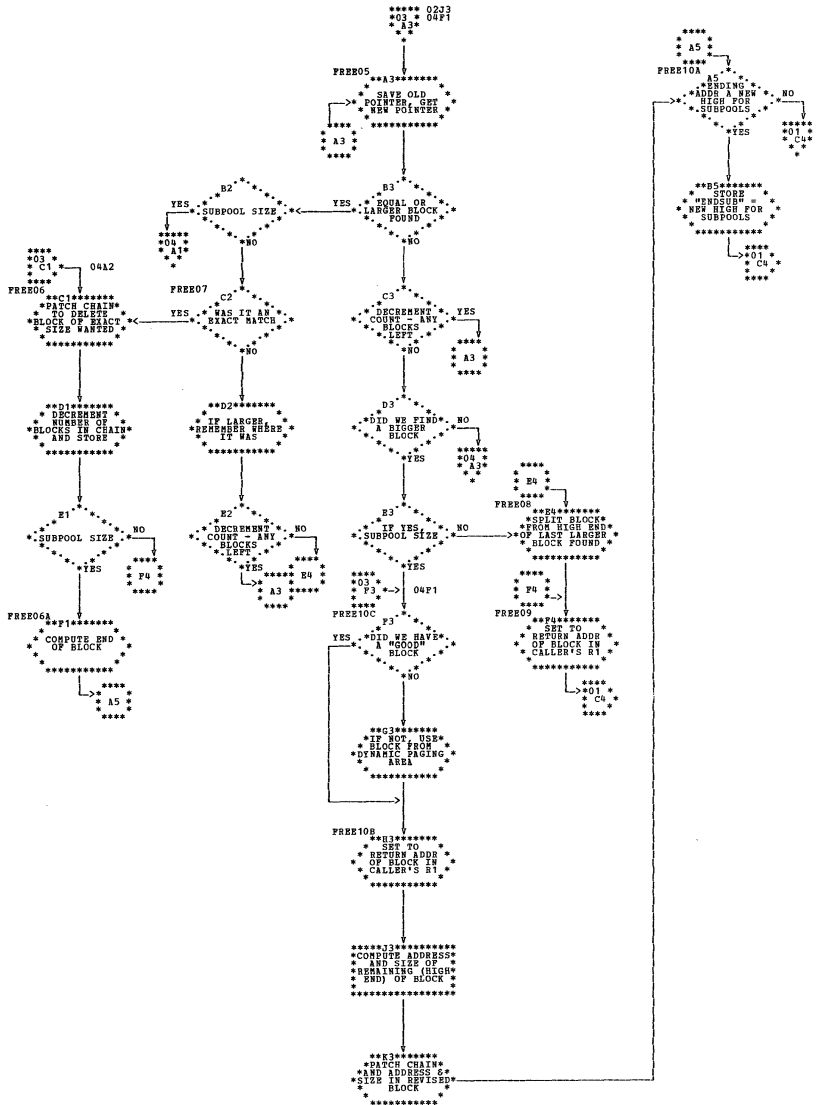
\* \* \* \* \*



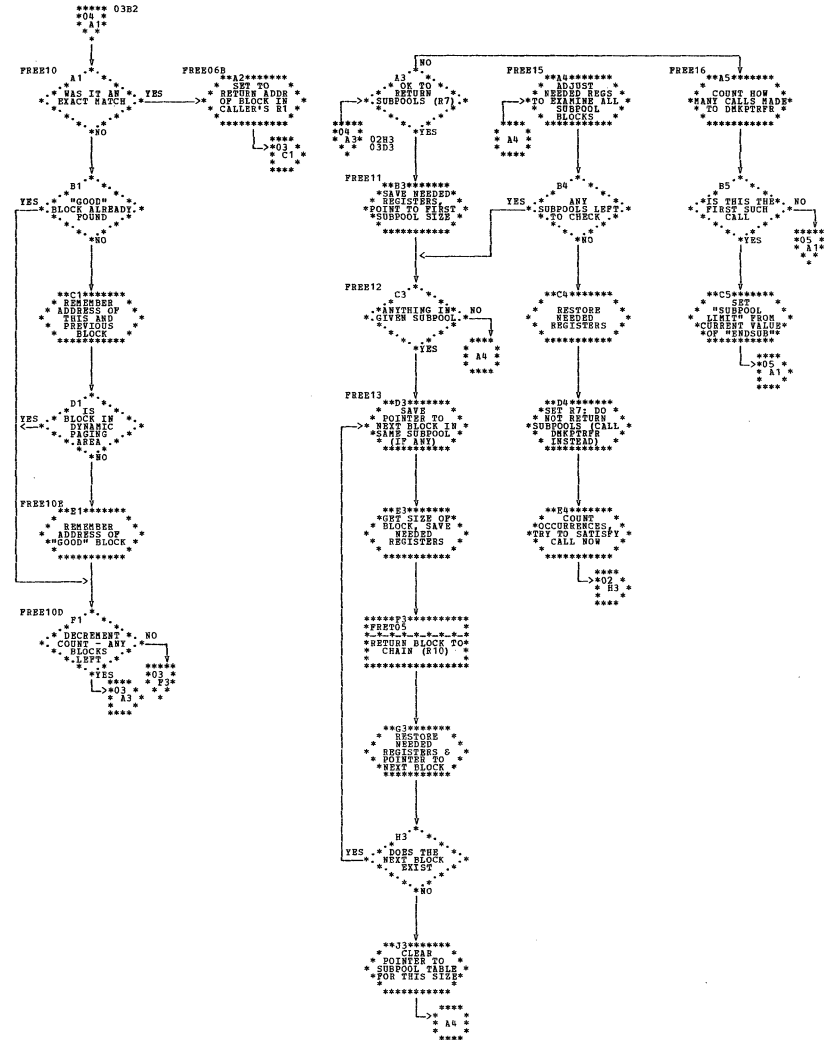
TO: P1  
07C1  
07D1  
08F2



| DMKFRE -- Free Storage Manager (Parts 3 and 4 of 9)

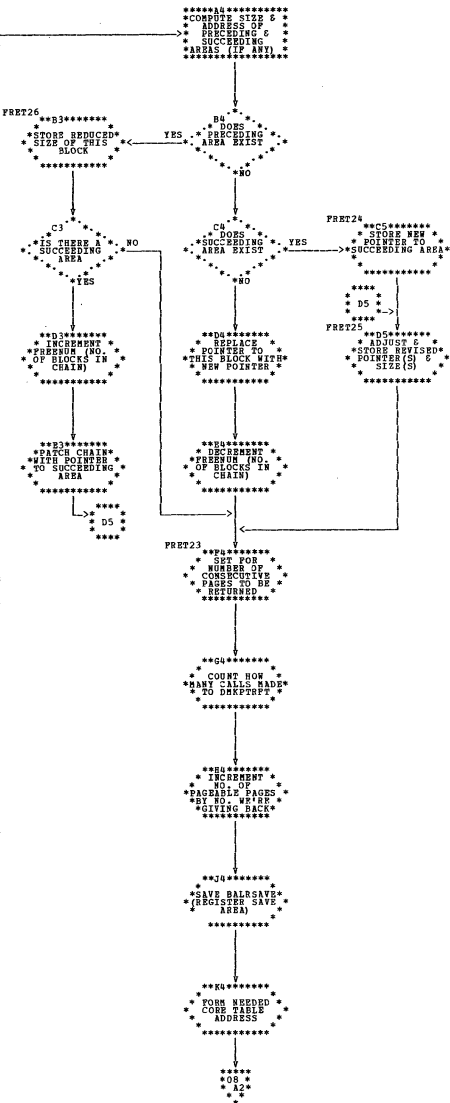
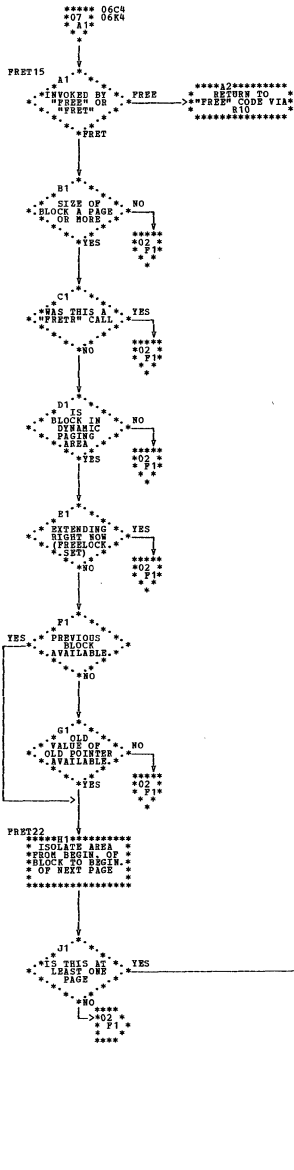


\* \* \* \* \*

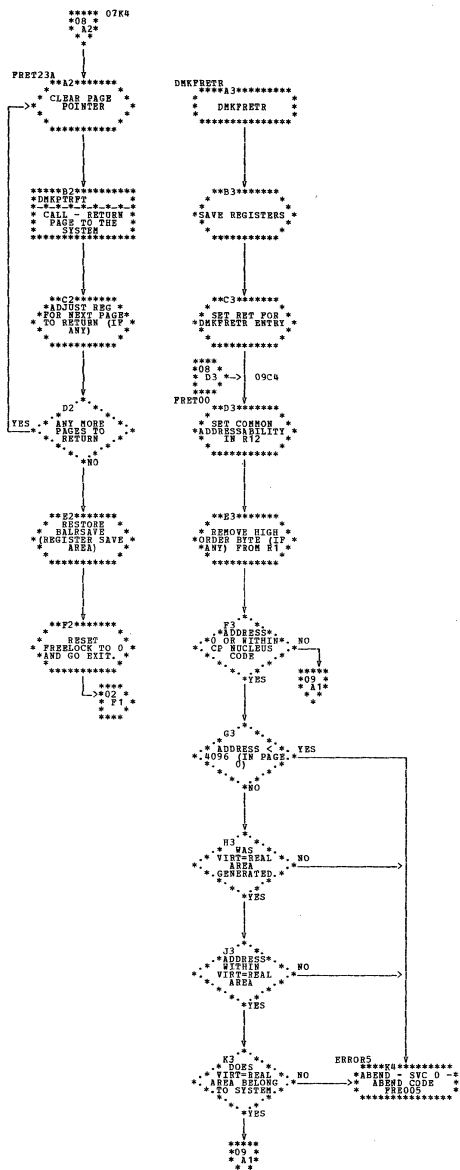




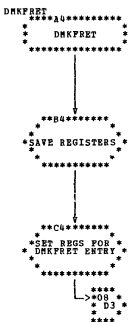
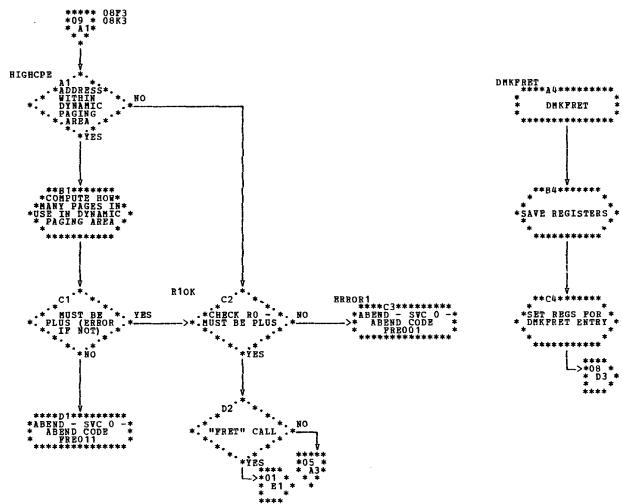
DMKFRE -- Free Storage Manager (Parts 7 and 8 of 9)



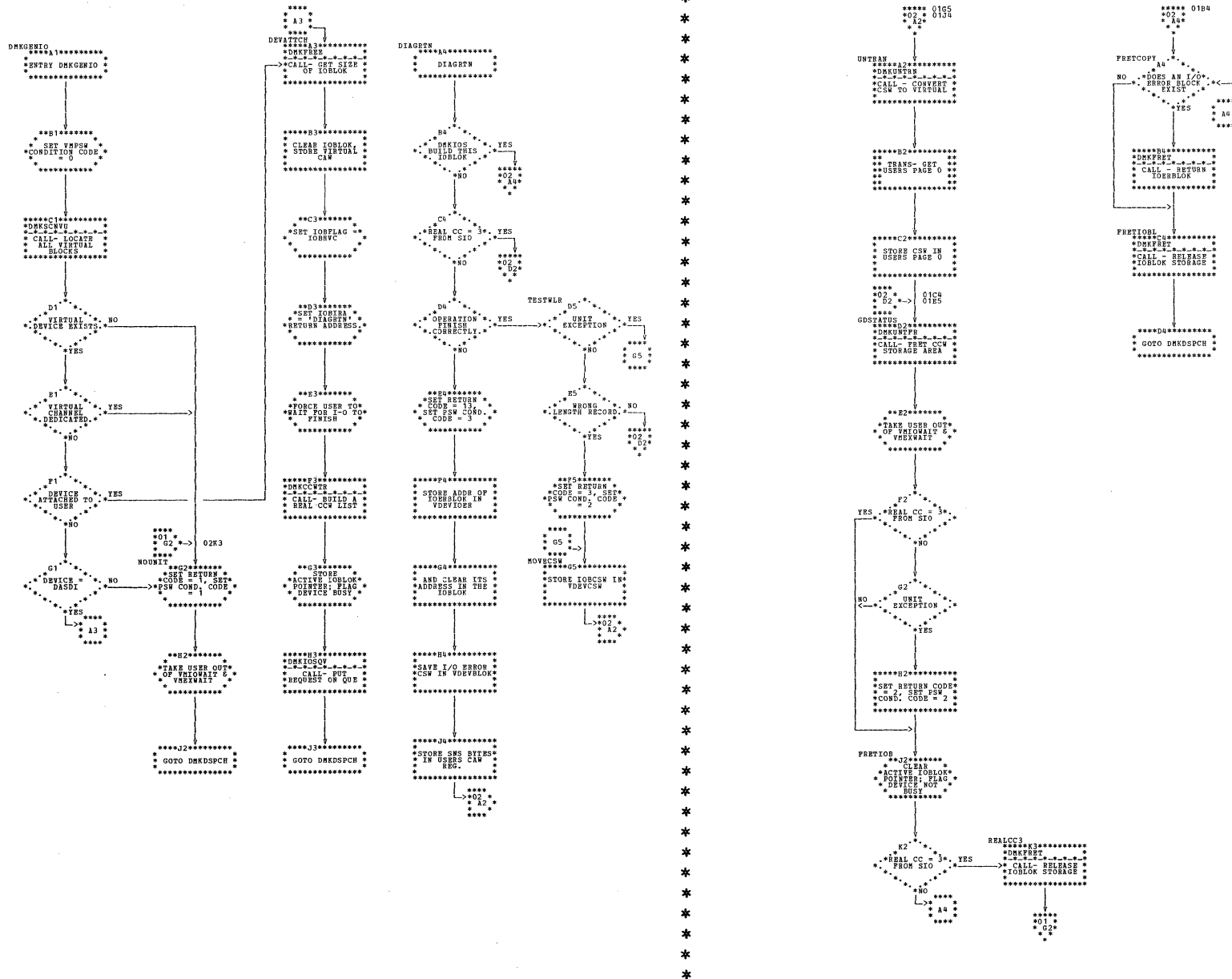
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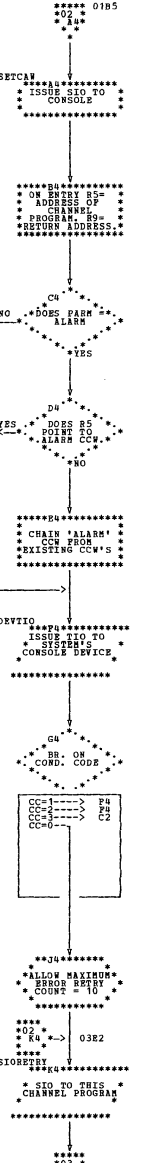
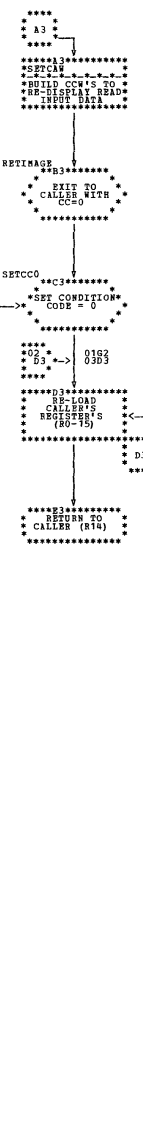
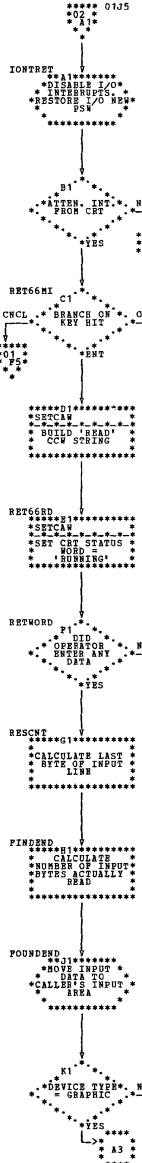
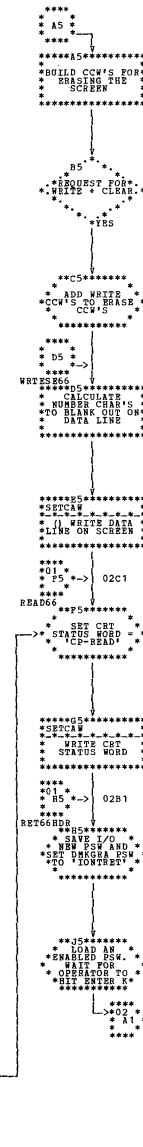
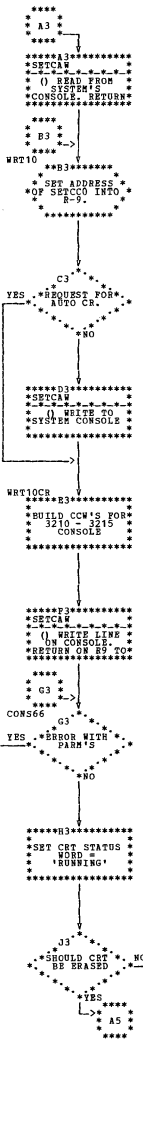
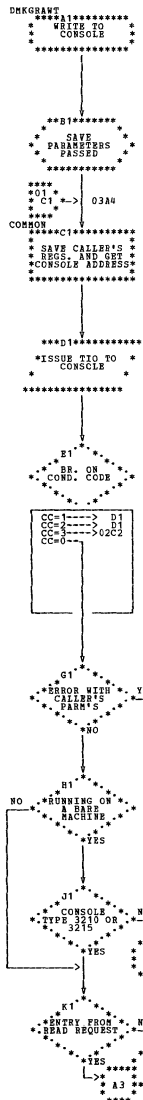


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ERRORS  
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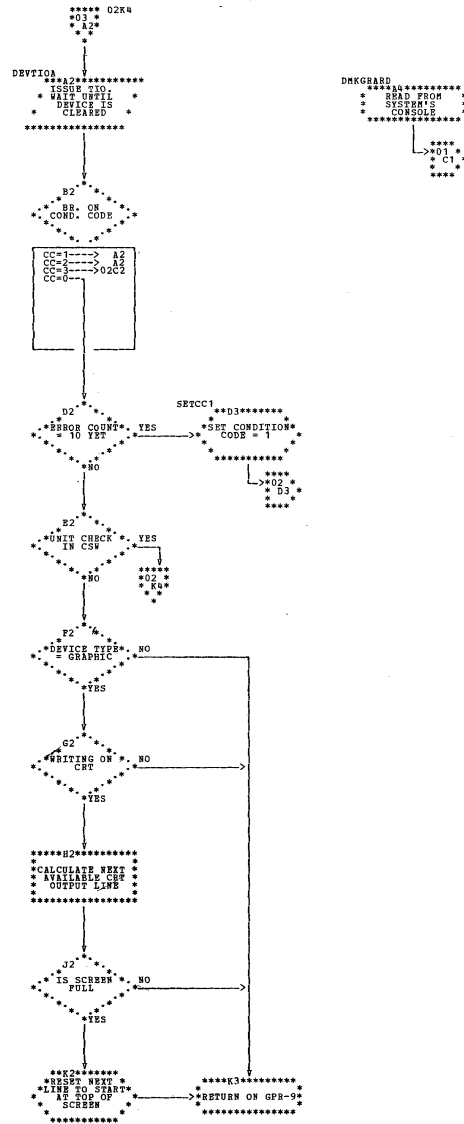
DMKGEN -- CMS I/O Error Recovery Interface (Parts 1 and 2 of 2)



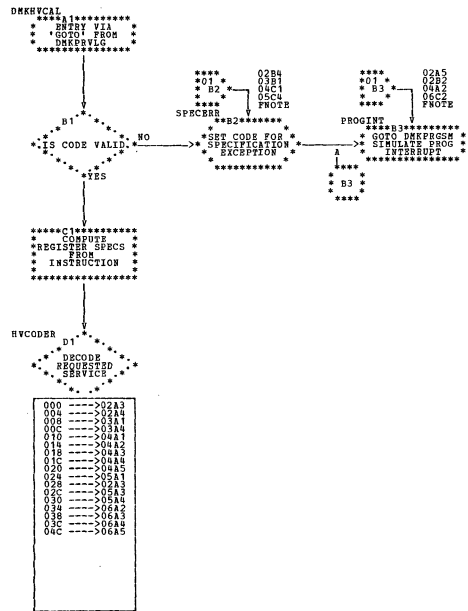


| DMKGRA -- System's Console Routine (Parts 1 and 2 of 3)

DMKGRA -- System's Console Routine (Part 3 of 3)



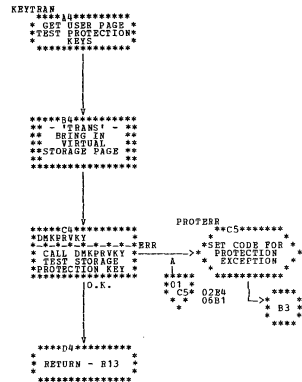




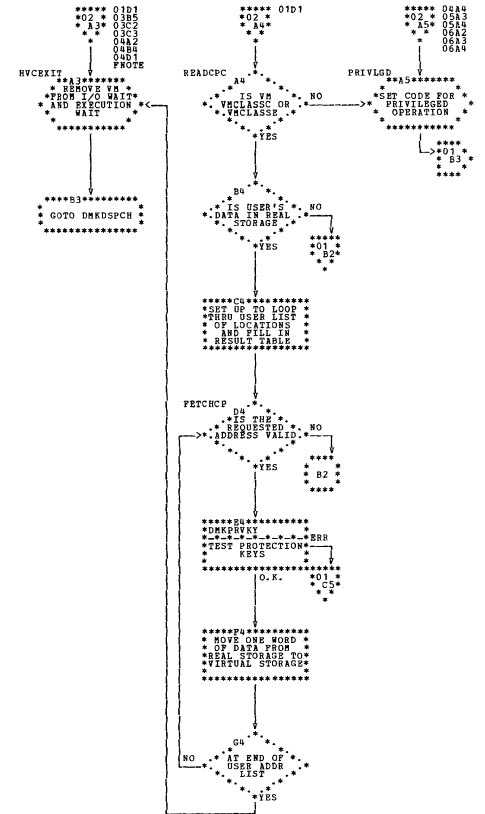
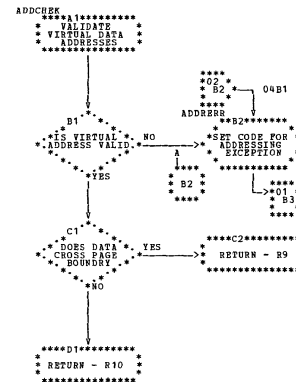
000	---	502A3
004	---	502B4
008	---	503A1
00C	---	503B4
010	---	504A1
014	---	504B4
018	---	504A3
01C	---	504B1
020	---	504A5
024	---	505A1
028	---	502A3
02C	---	502A3
030	---	505A4
034	---	506A2
038	---	506A3
03C	---	506A8
040	---	506A5

TO: B2  
02B4  
06C5

TO: B3  
02B3  
06C4



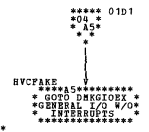
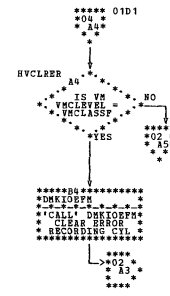
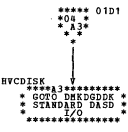
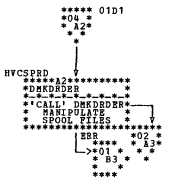
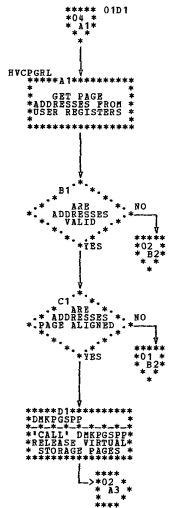
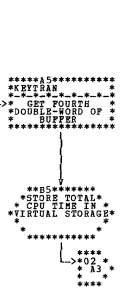
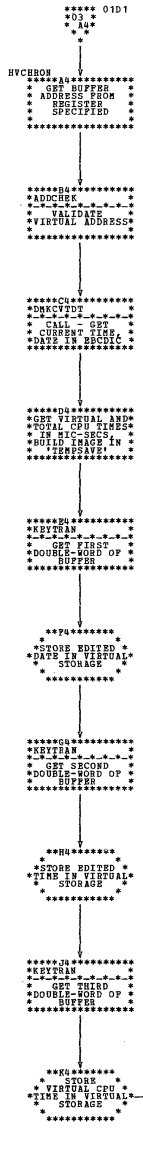
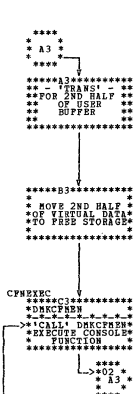
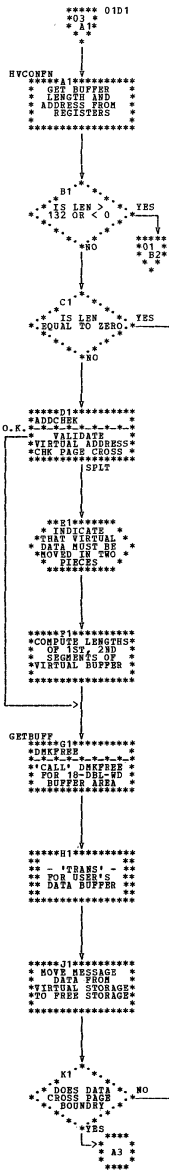
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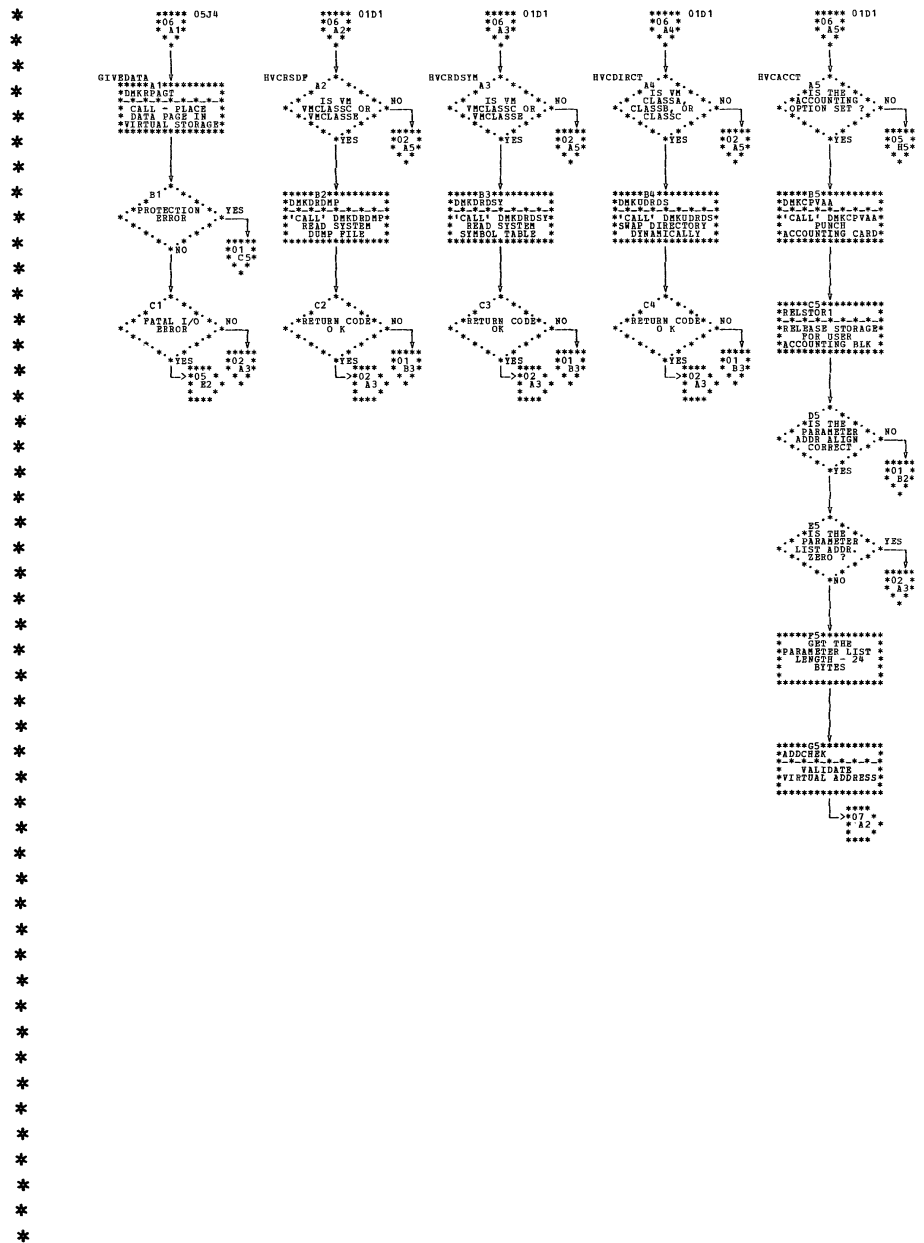
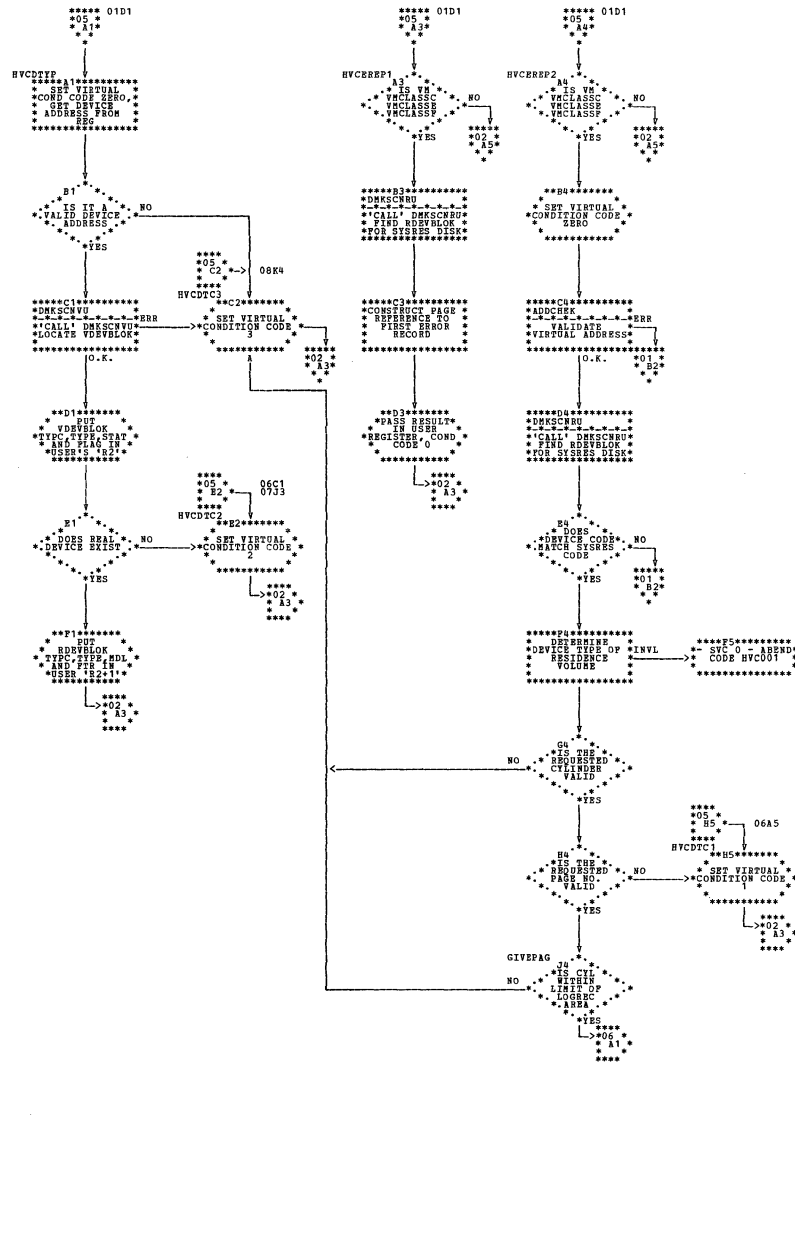
TO: A3  
05C2  
05D  
05E  
05F  
05H  
05I  
05J  
05K  
05L  
05M  
05N  
05O

| DMKHVC -- Process DIAGNOSE Instructions (Parts 1 and 2 of 8)

| DMKHVC -- Process DIAGNOSE Instructions (Parts 3 and 4 of 8)

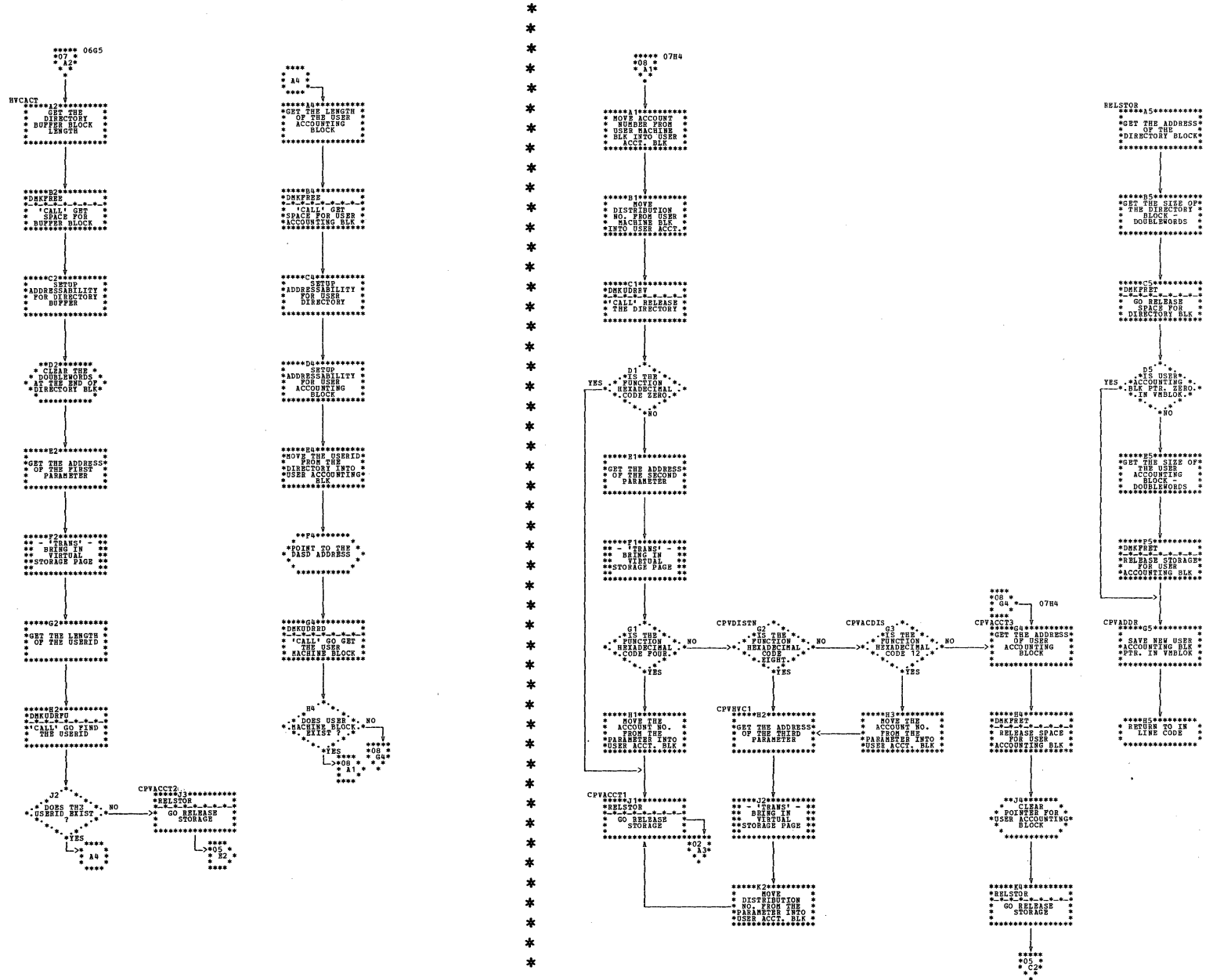


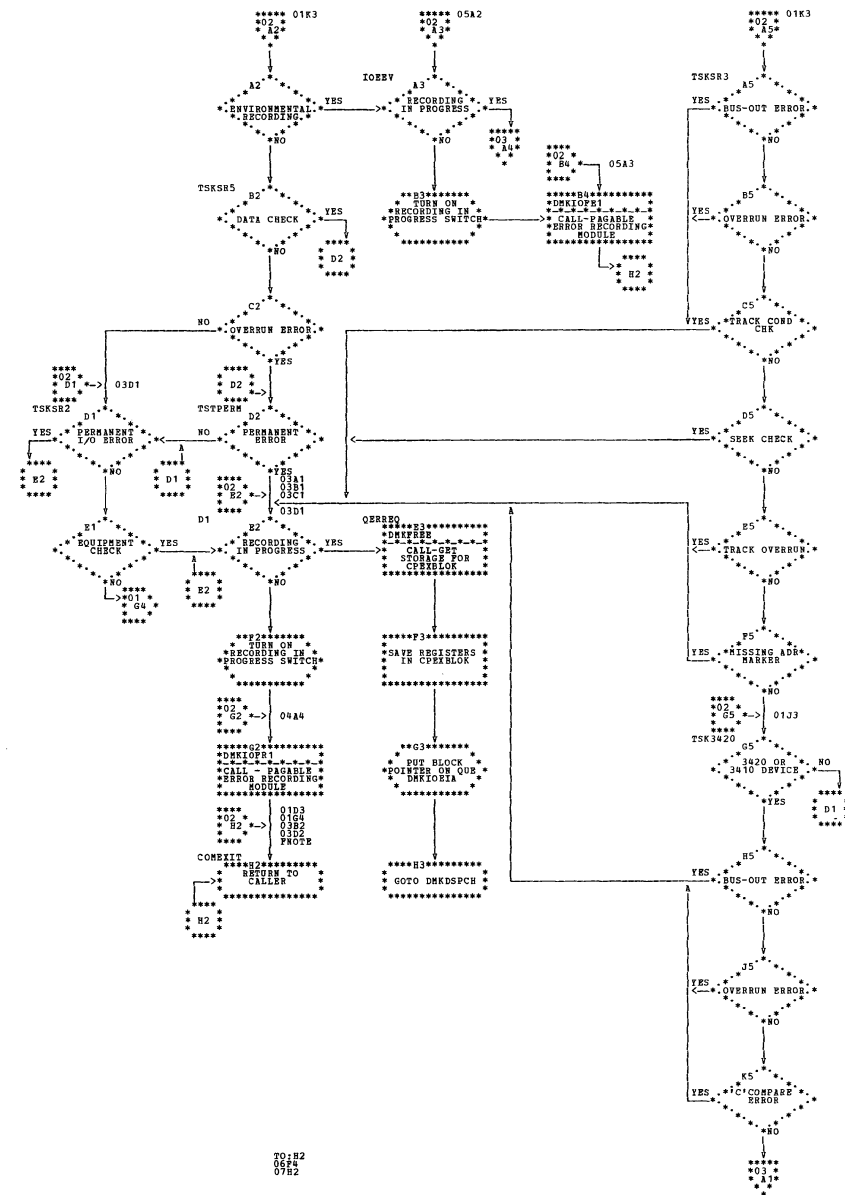
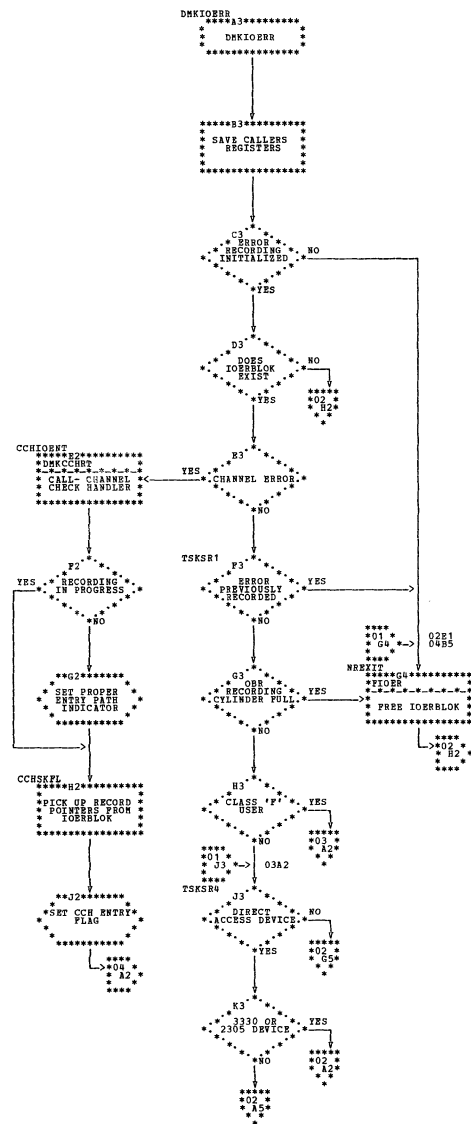
\* \* \* \* \*



| DMKHVC -- Process DIAGNOSE Instructions (Parts 5 and 6 of 8)

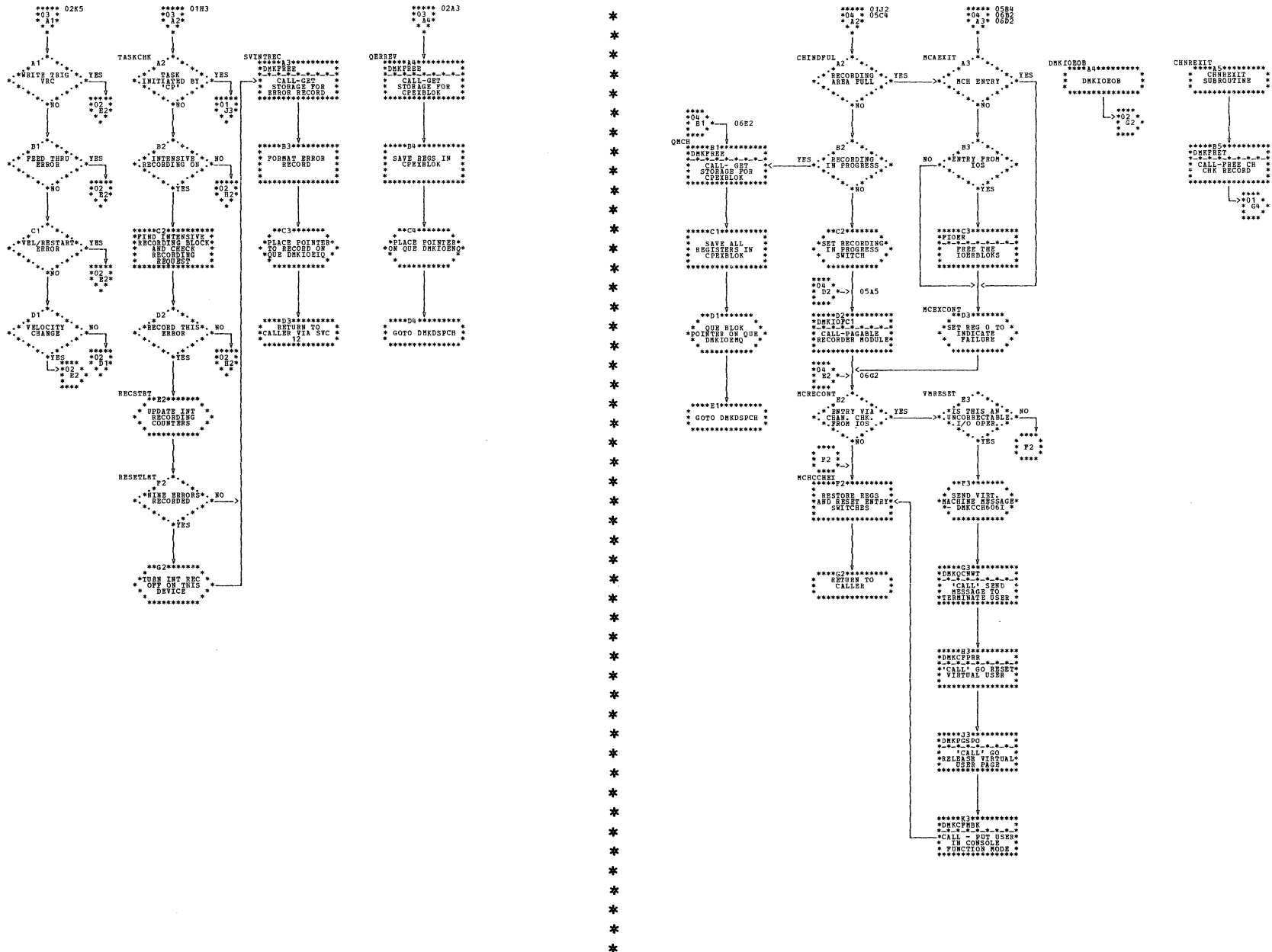
DMKHVC -- Process DIAGNOSE Instructions (Parts 7 and 8 of 8)





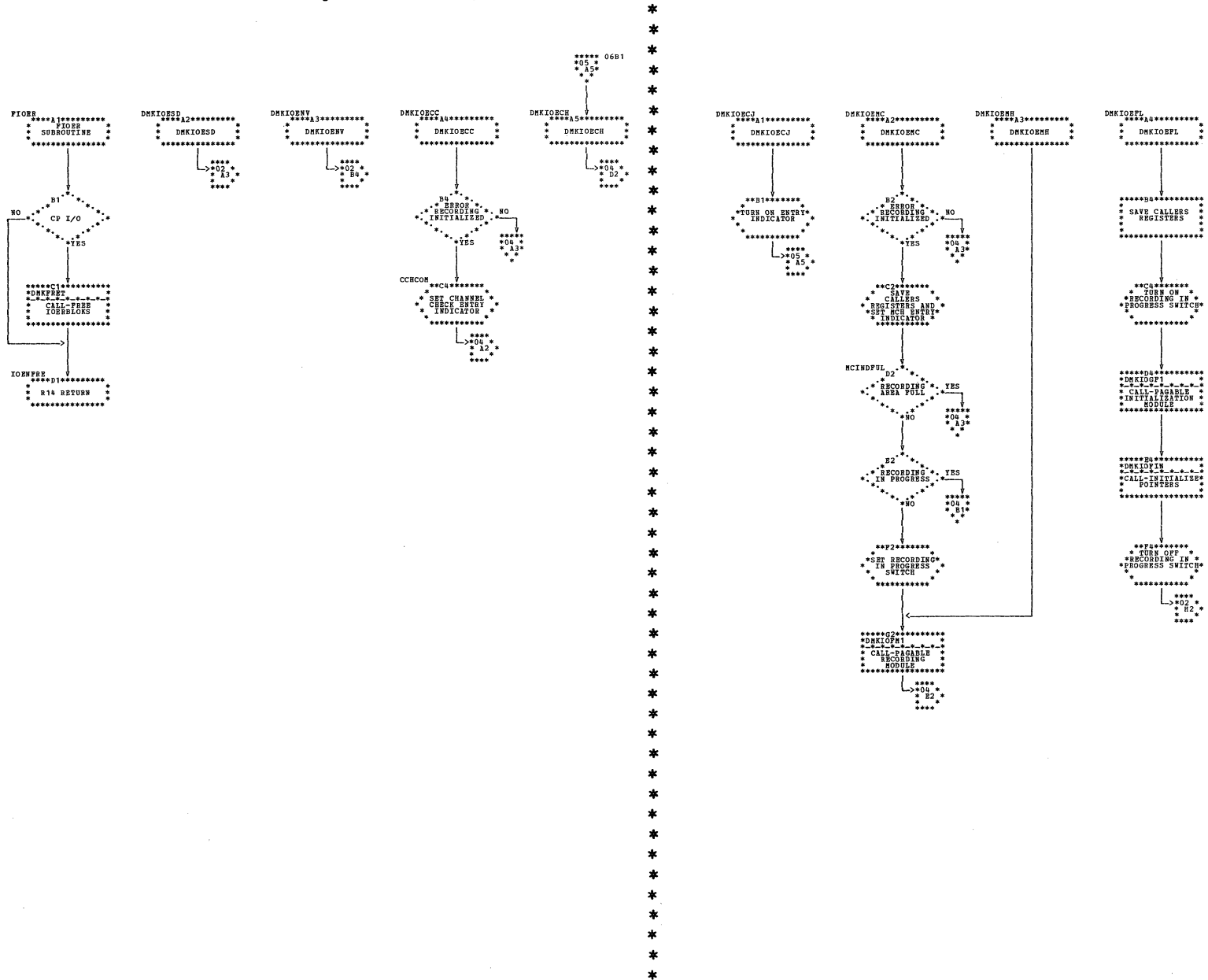
| DMKIOE -- Main Error Recording Processor (Parts 1 and 2 of 7)



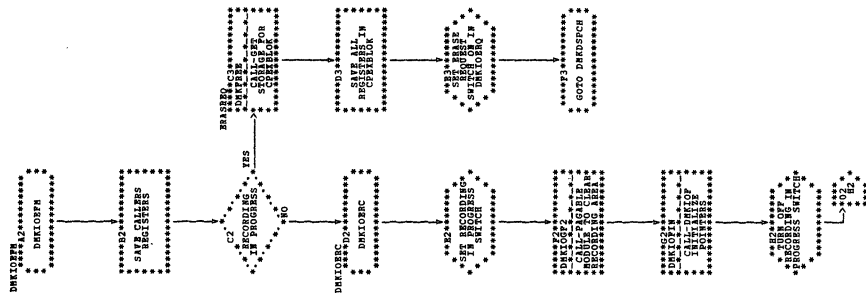


| DMKIOE -- Main Error Recording Processor (Parts 3 and 4 of 7)

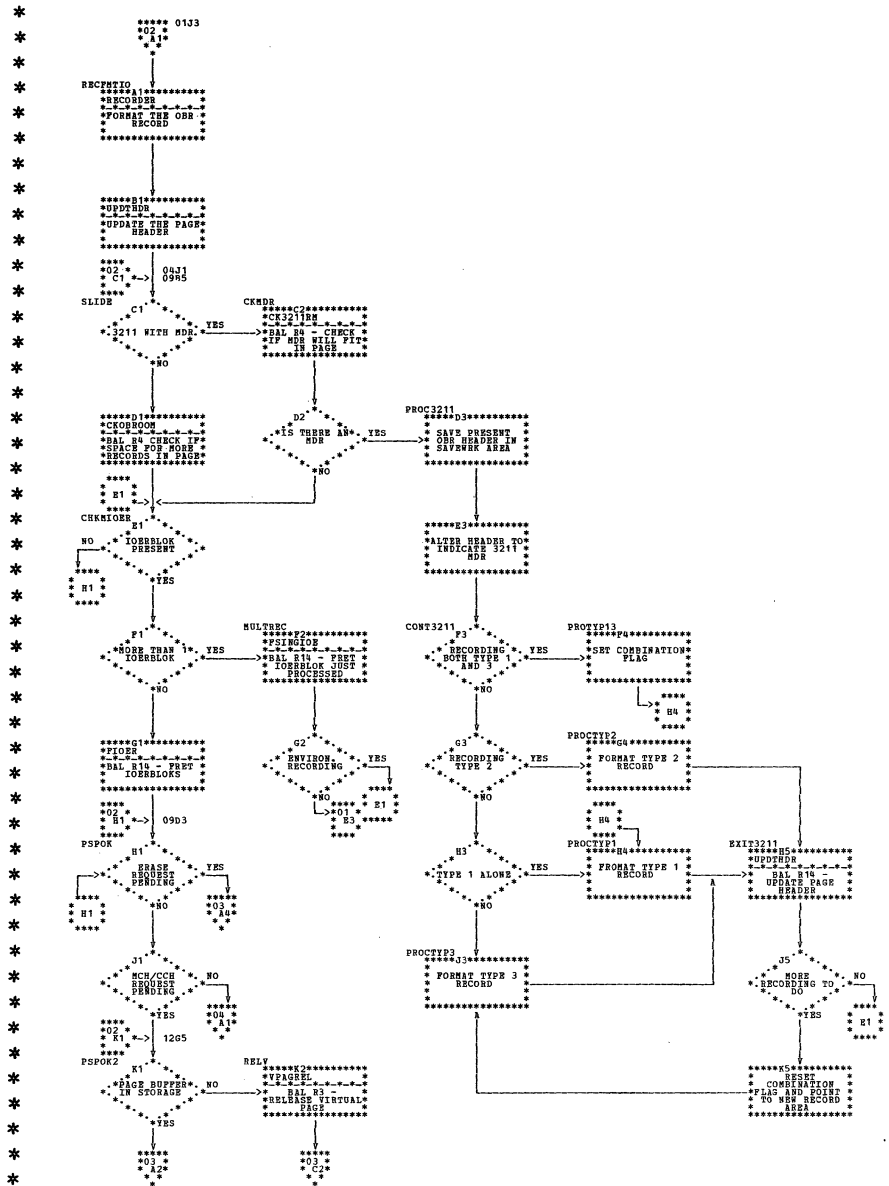
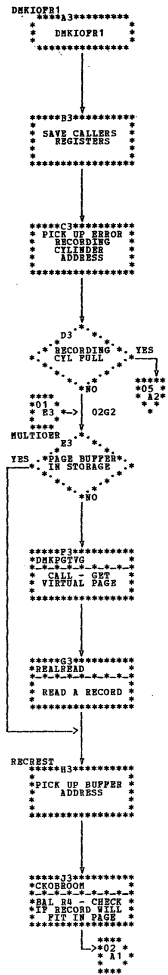
| DMKIOE -- Main Error Recording Processor (Parts 5 and 6 of 7)







| DMKIOF -- Error Recorder (Parts 1 and 2 of 13)

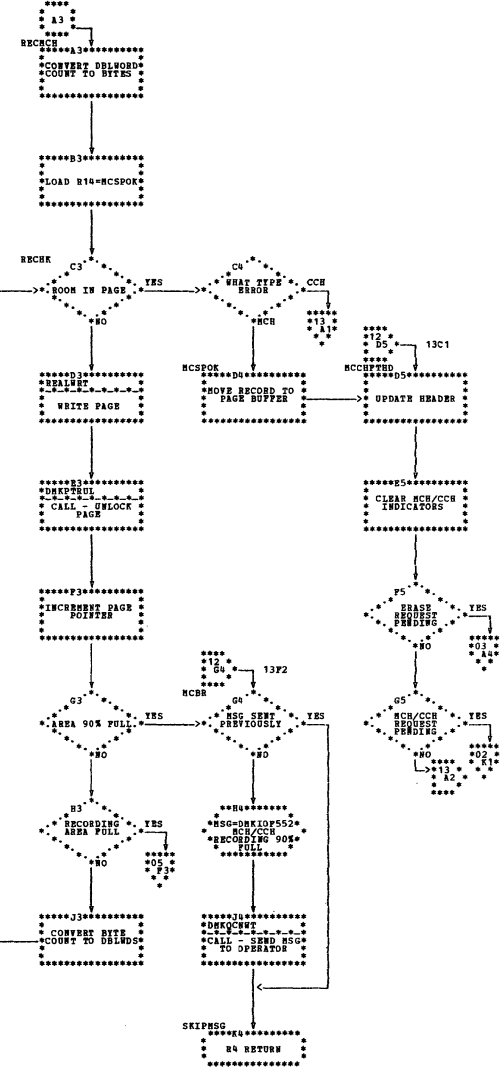
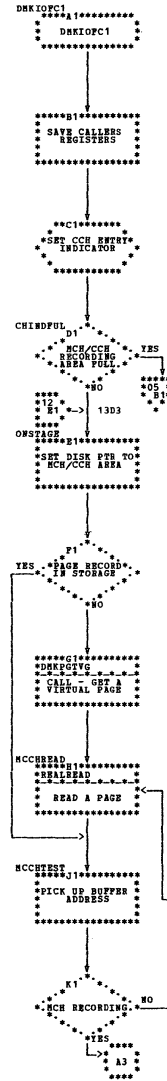
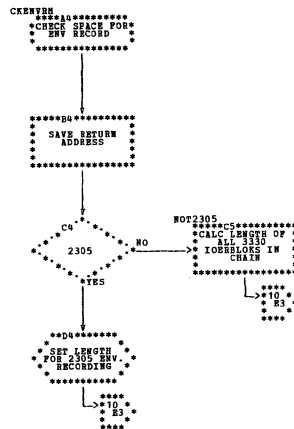
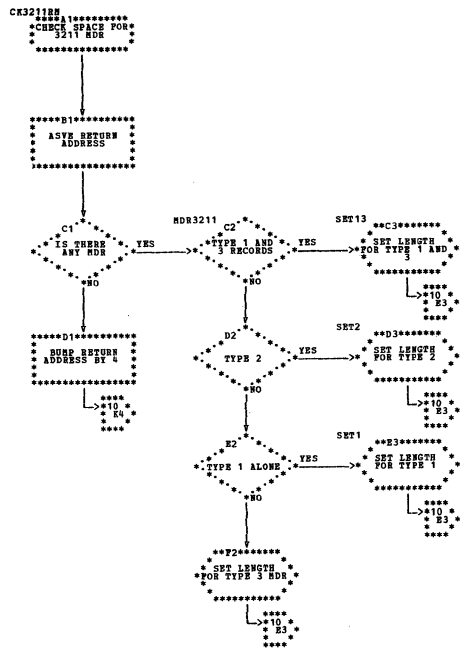




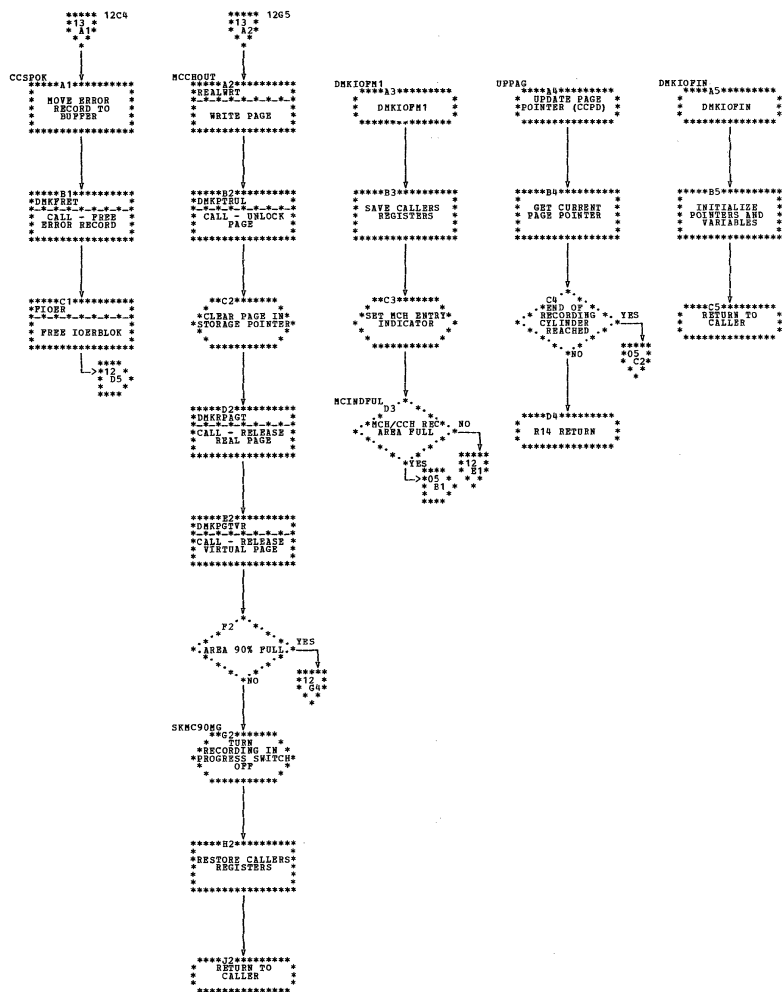




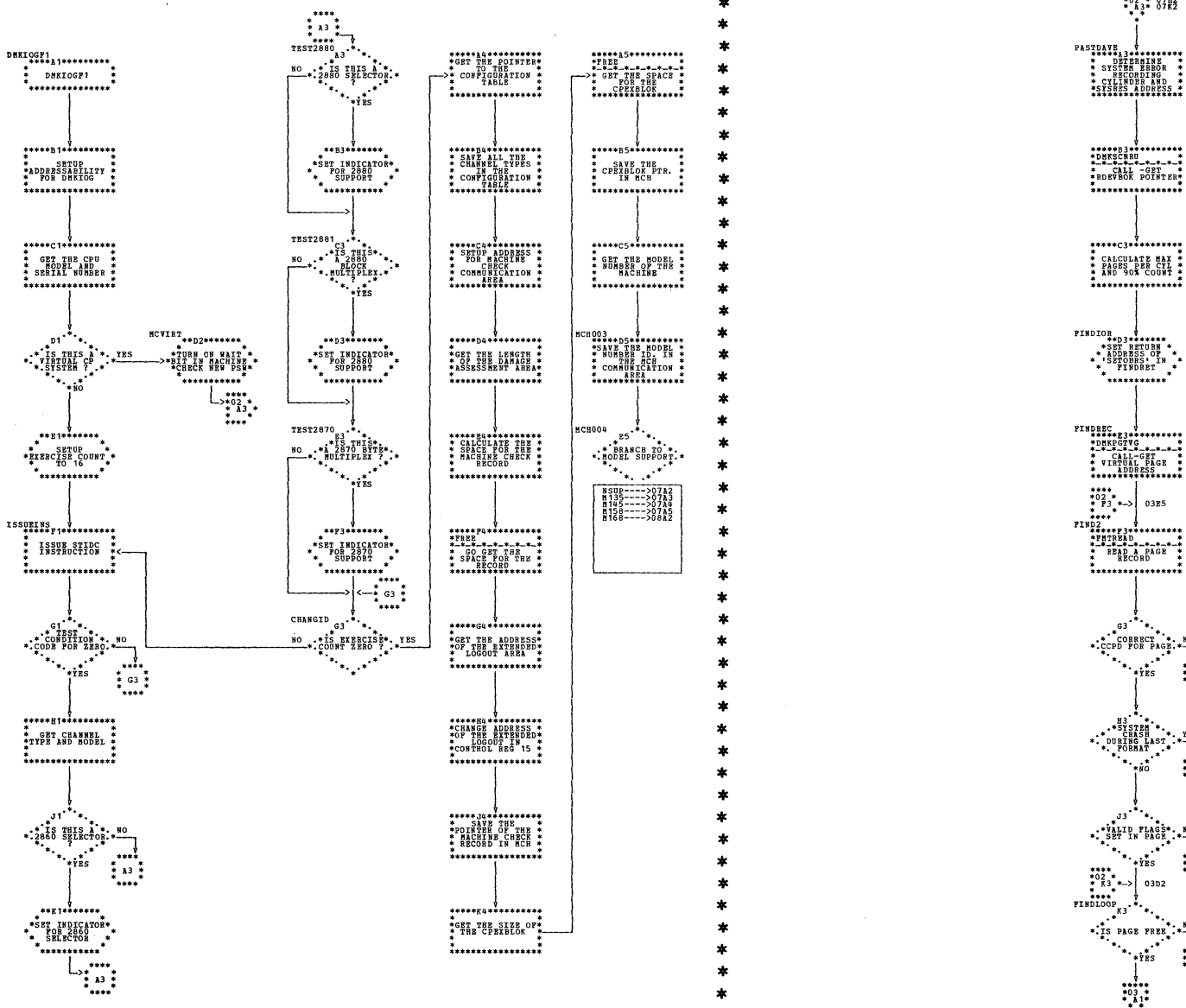




| DMKIOF -- Error Recorder (Parts 13 of 13)

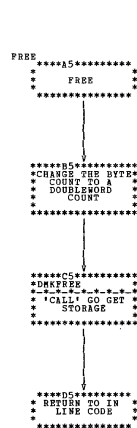
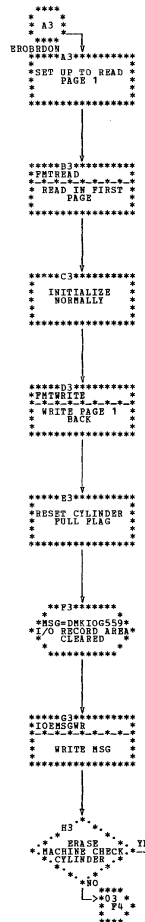
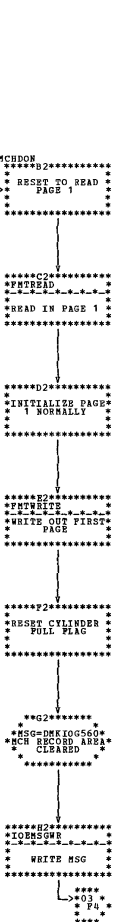
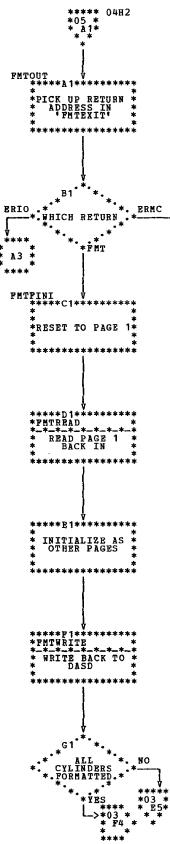




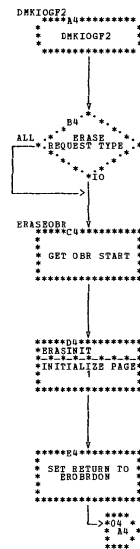
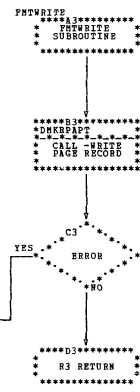
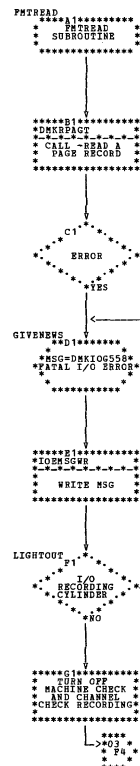


| DMKIOG -- Error Recording, Initialization and Clearing (Parts 1 and 2 of 8)



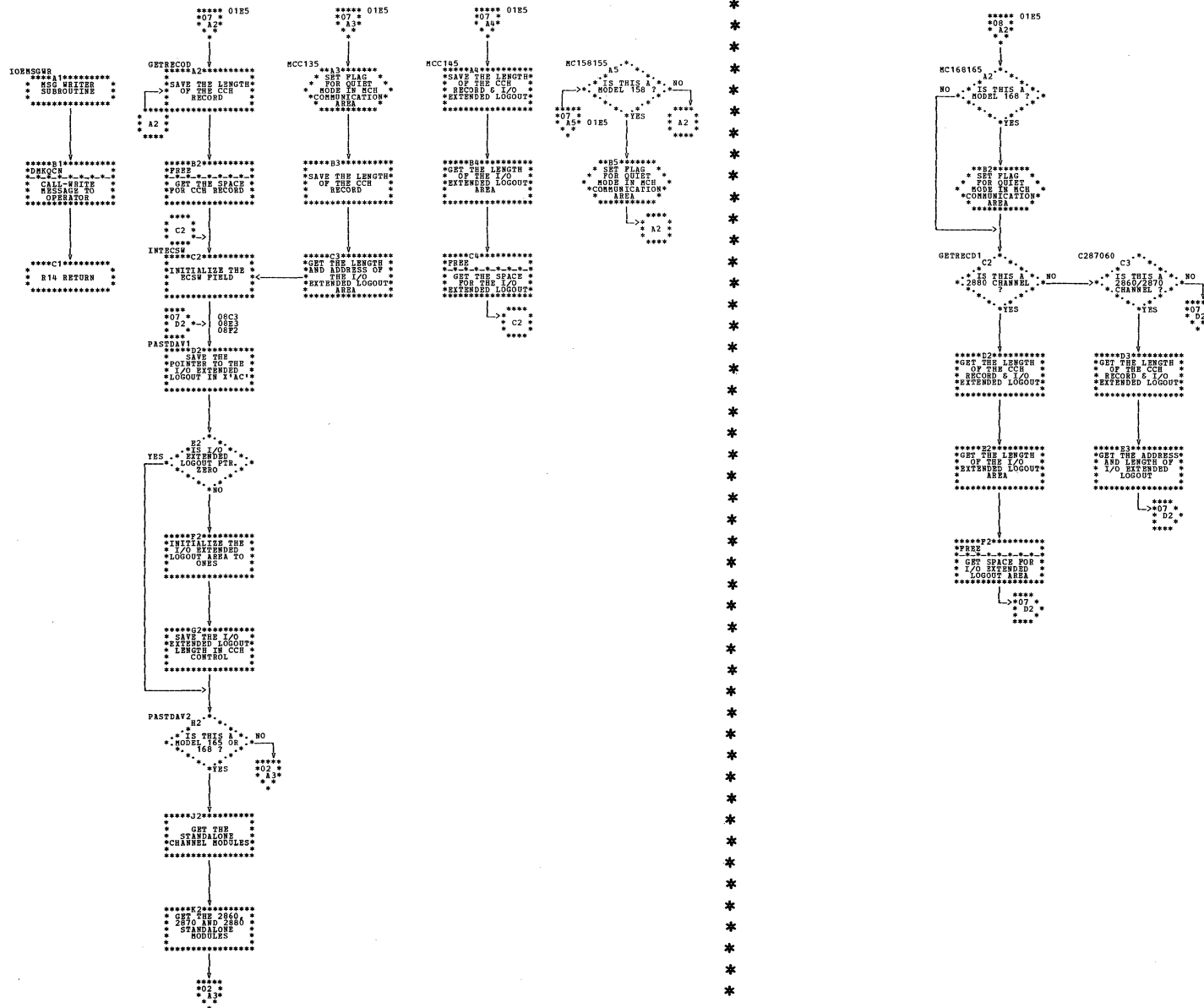


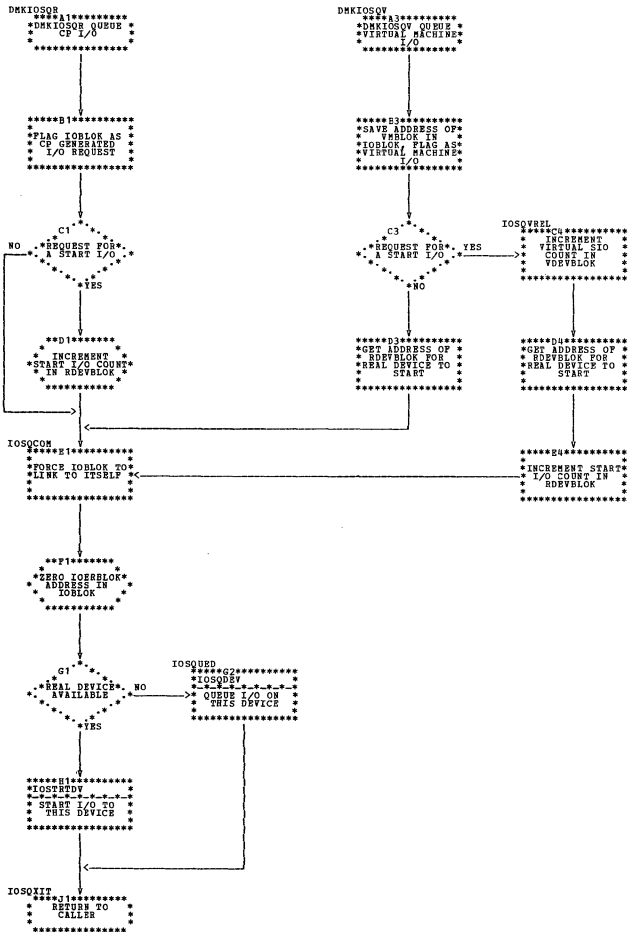
\* \* \* \* \*



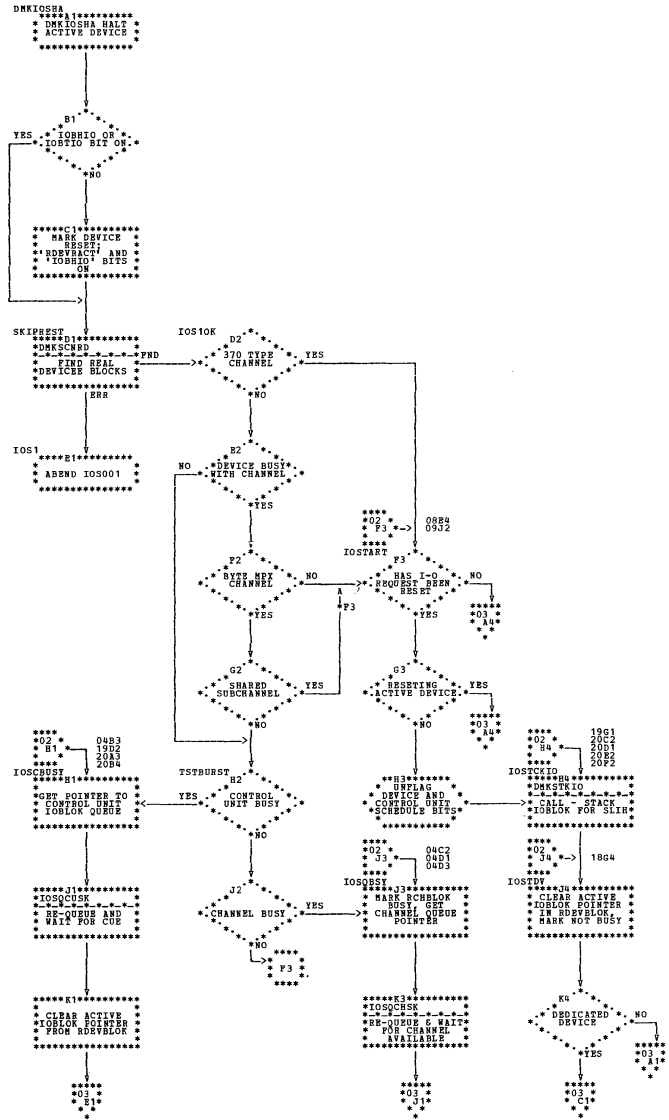
| DMKIOG -- Error Recording, Initialization and Clearing (Parts 5 and 6 of 8)

DMKIOG -- Error Recording, Initialization and Clearing (Parts 7 and 8 of 8)





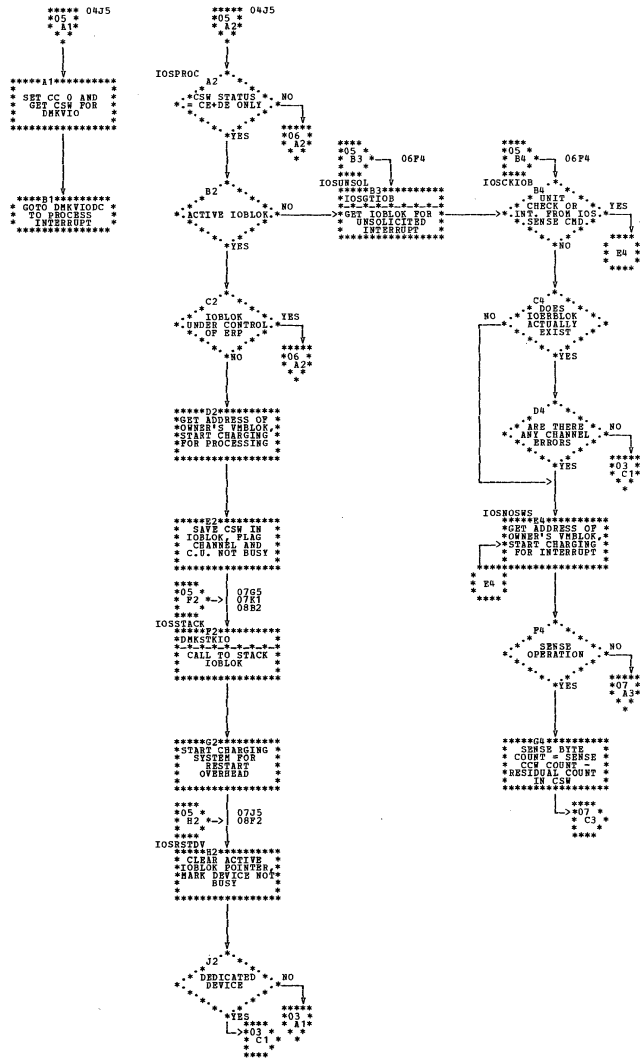
| DMKIOS -- I/O Supervisor (Parts 1 and 2 of 20)



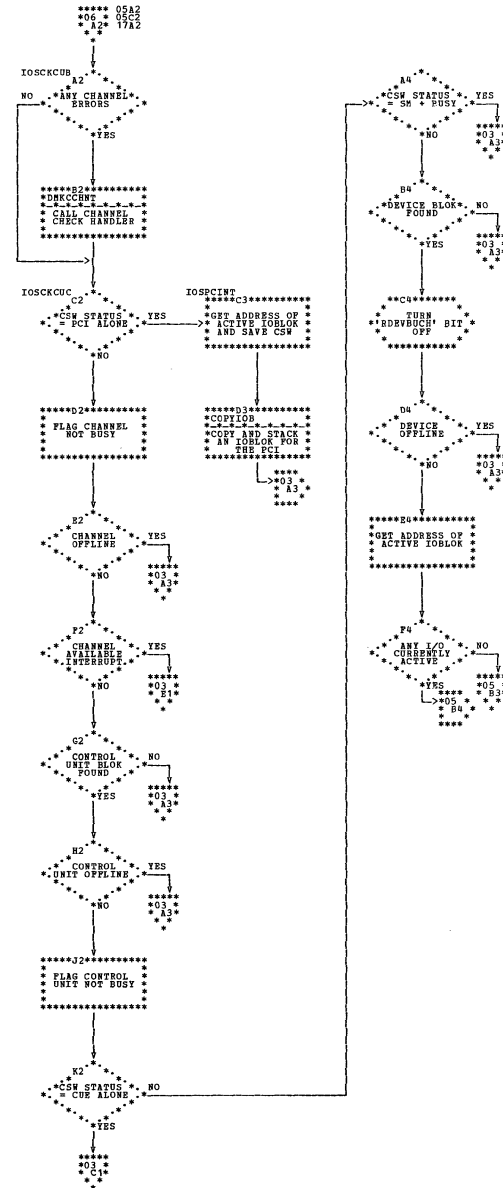




DMKIOS -- I/O Supervisor (Parts 5 and 6 of 20)



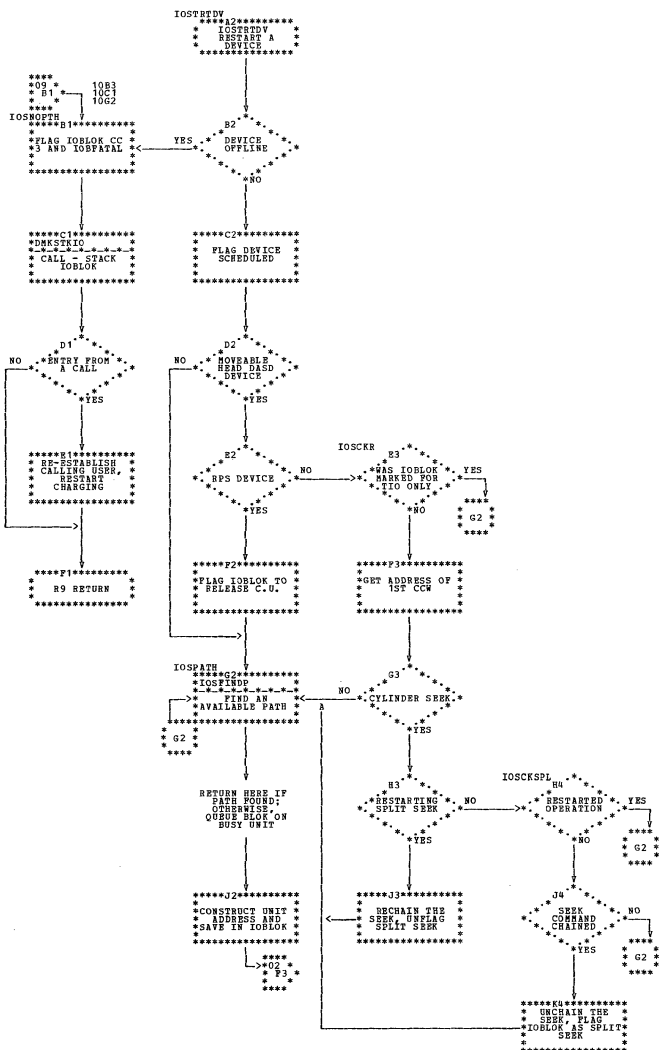
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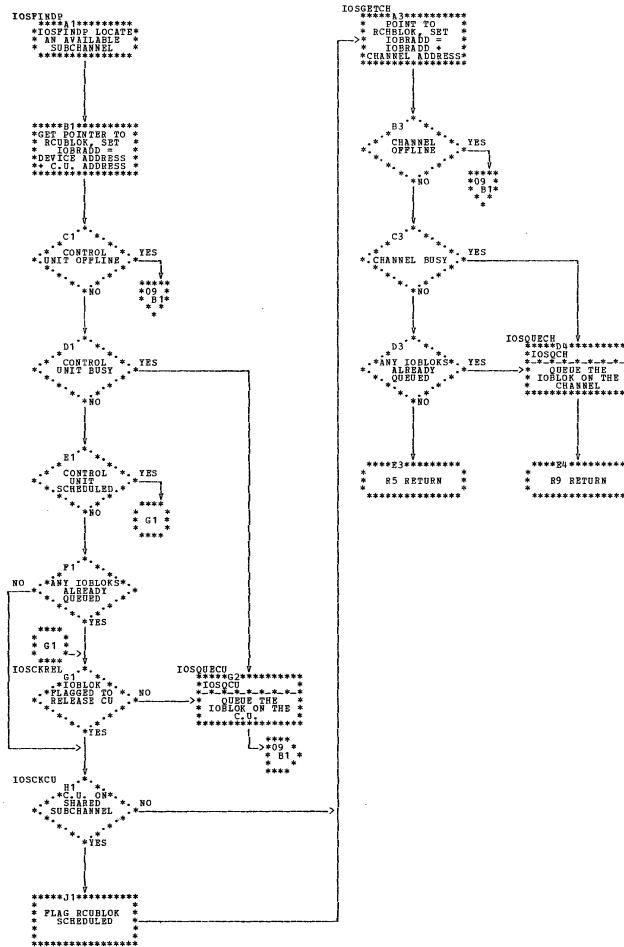


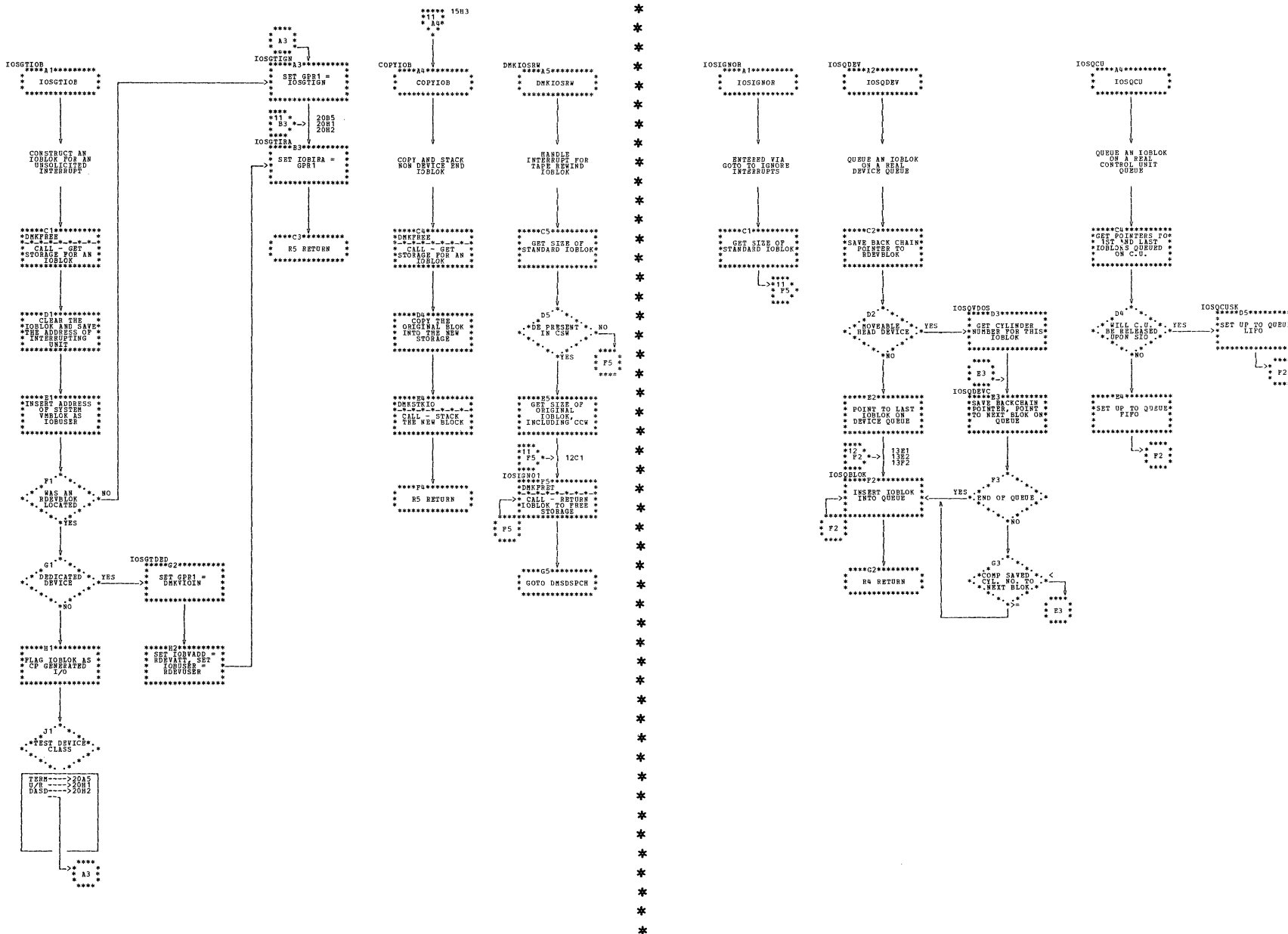


DMKIOS -- I/O Supervisor (Parts 9 and 10 of 20)



\* \* \* \* \*



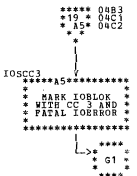
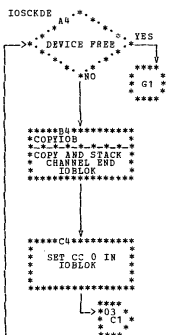
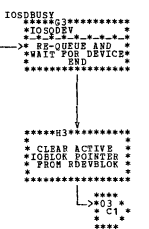
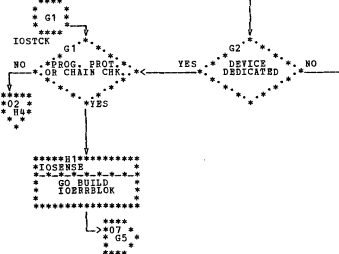
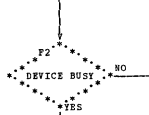
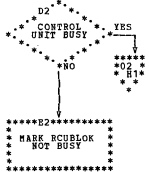
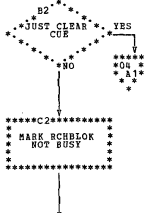
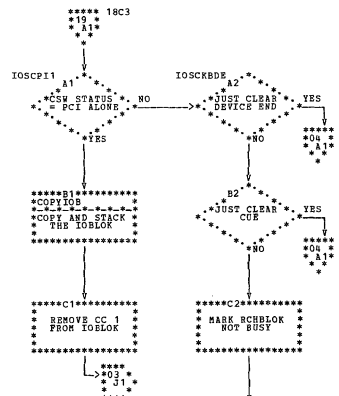


| DMKIOS -- I/O Supervisor (Parts 11 and 12 of 20)

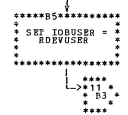
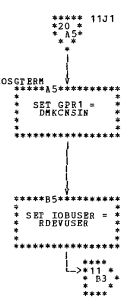
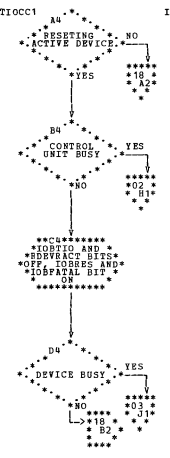
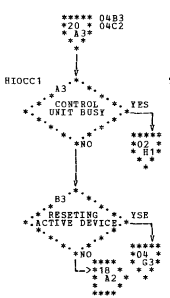
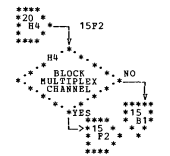
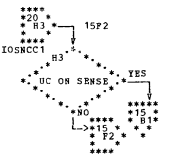
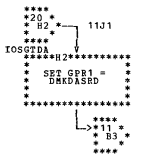
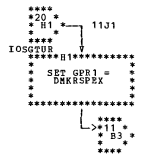
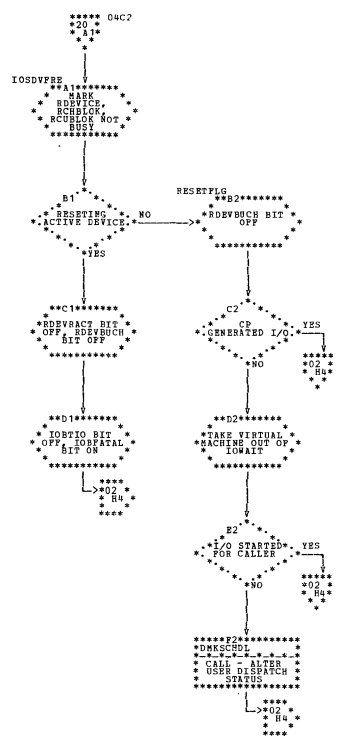




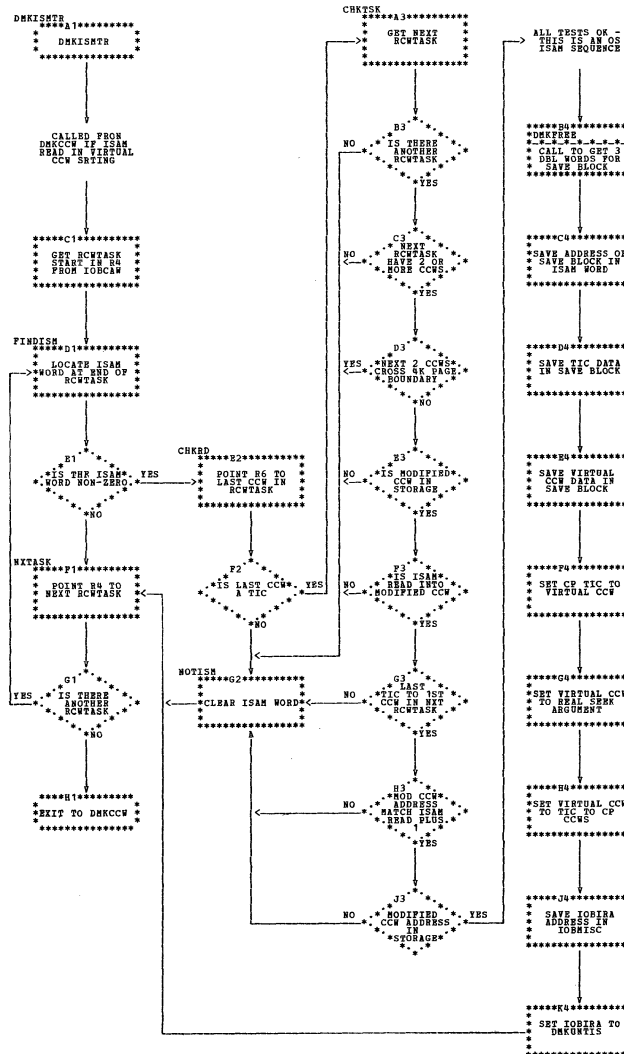




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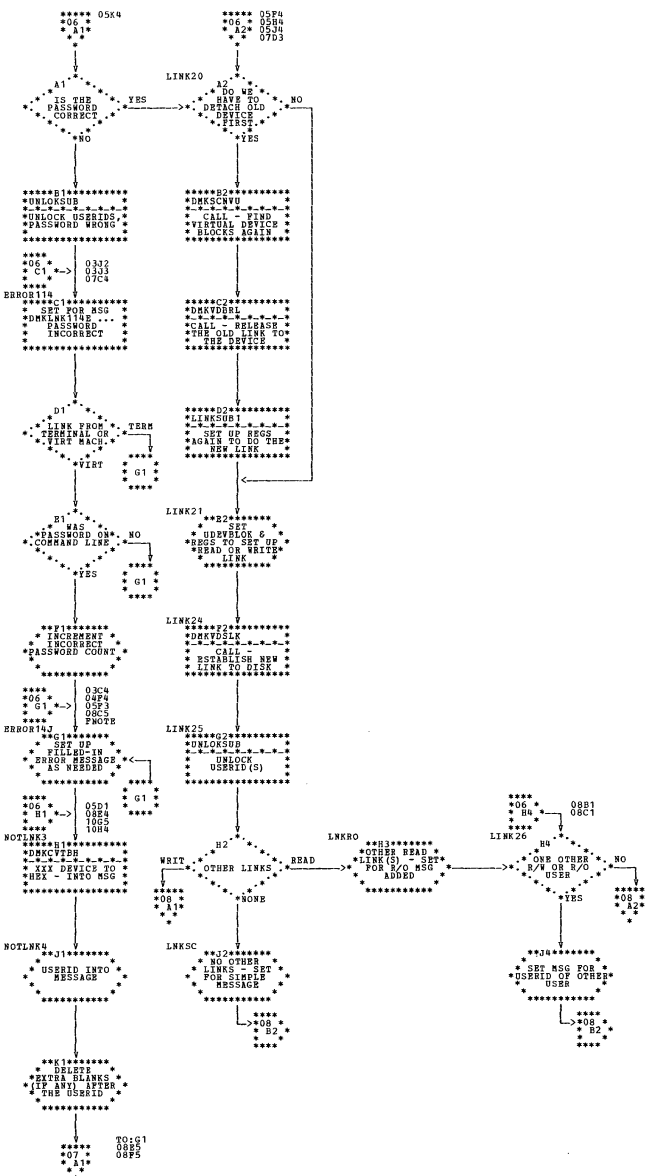
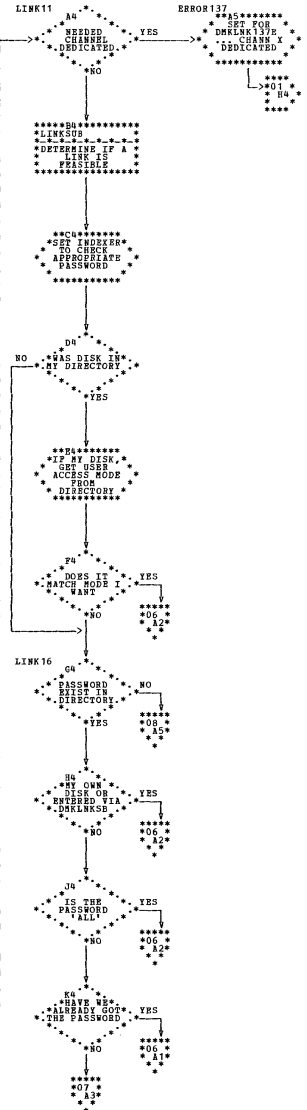
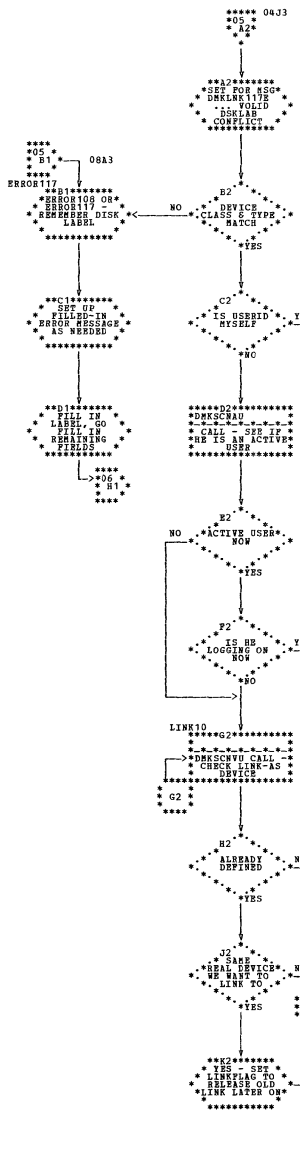
DMKISM -- Modify RCWTASK for OS ISAM Input/Output (Part 1 of 1)





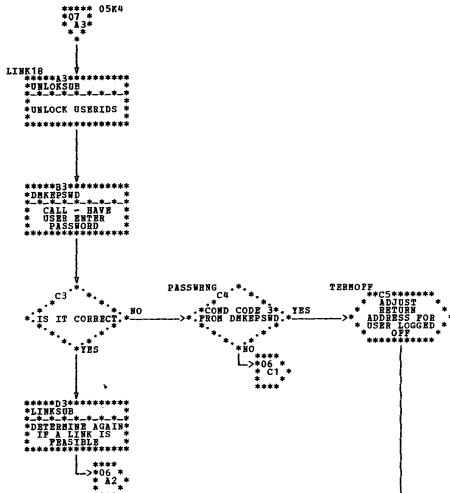
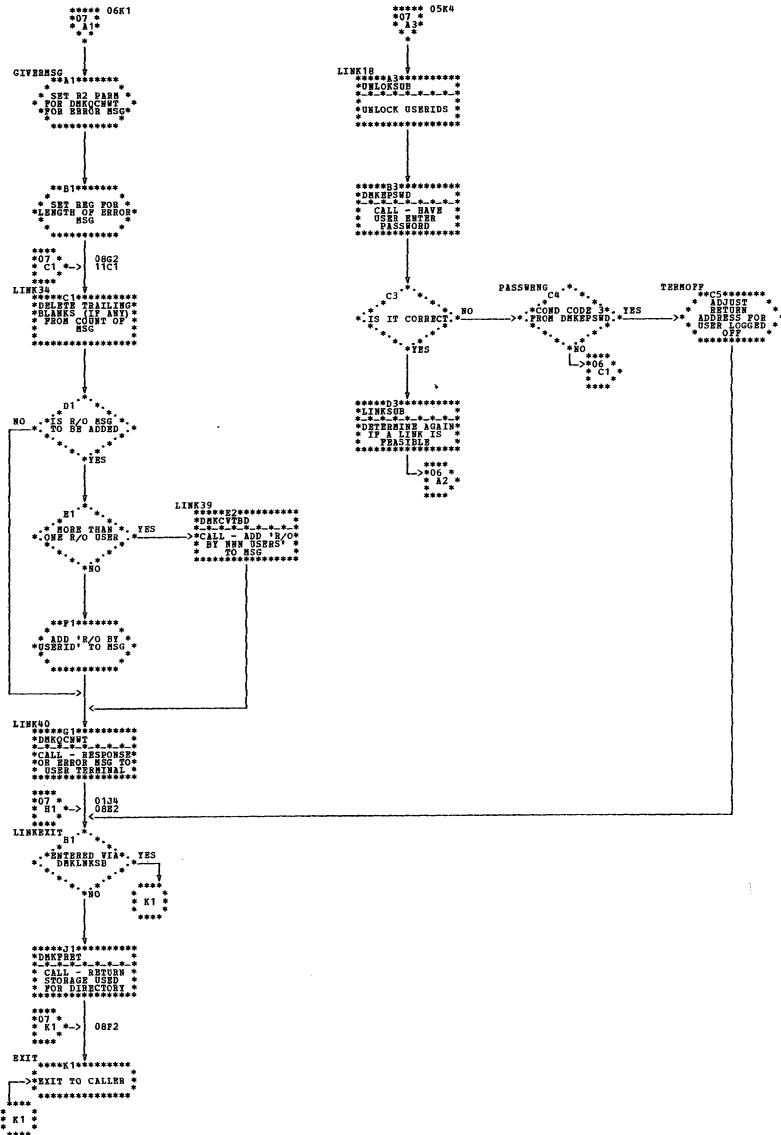




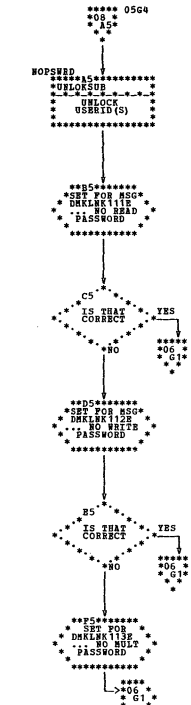
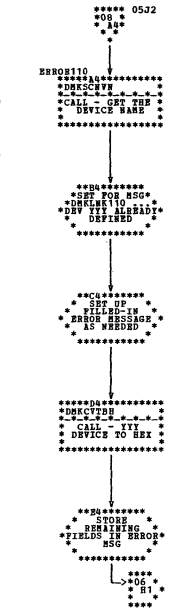
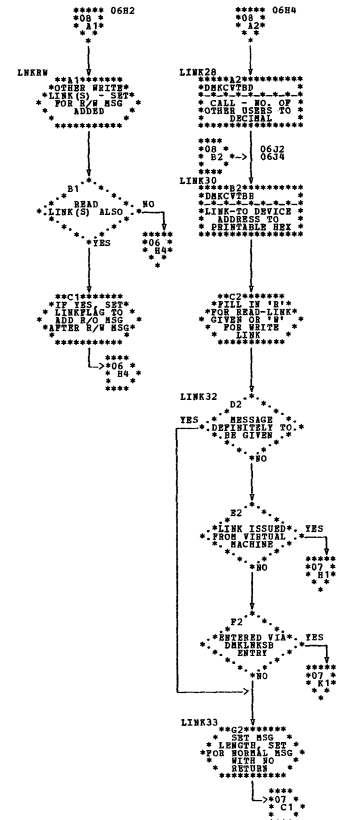


DMKLNK -- Process LINK Command; Link to a Virtual DASD (Parts 5 and 6 of 11)

DMKLNK -- Process LINK Command; Link to a Virtual DASD (Parts 7 and 8 of 11)



\* \* \* \* \*



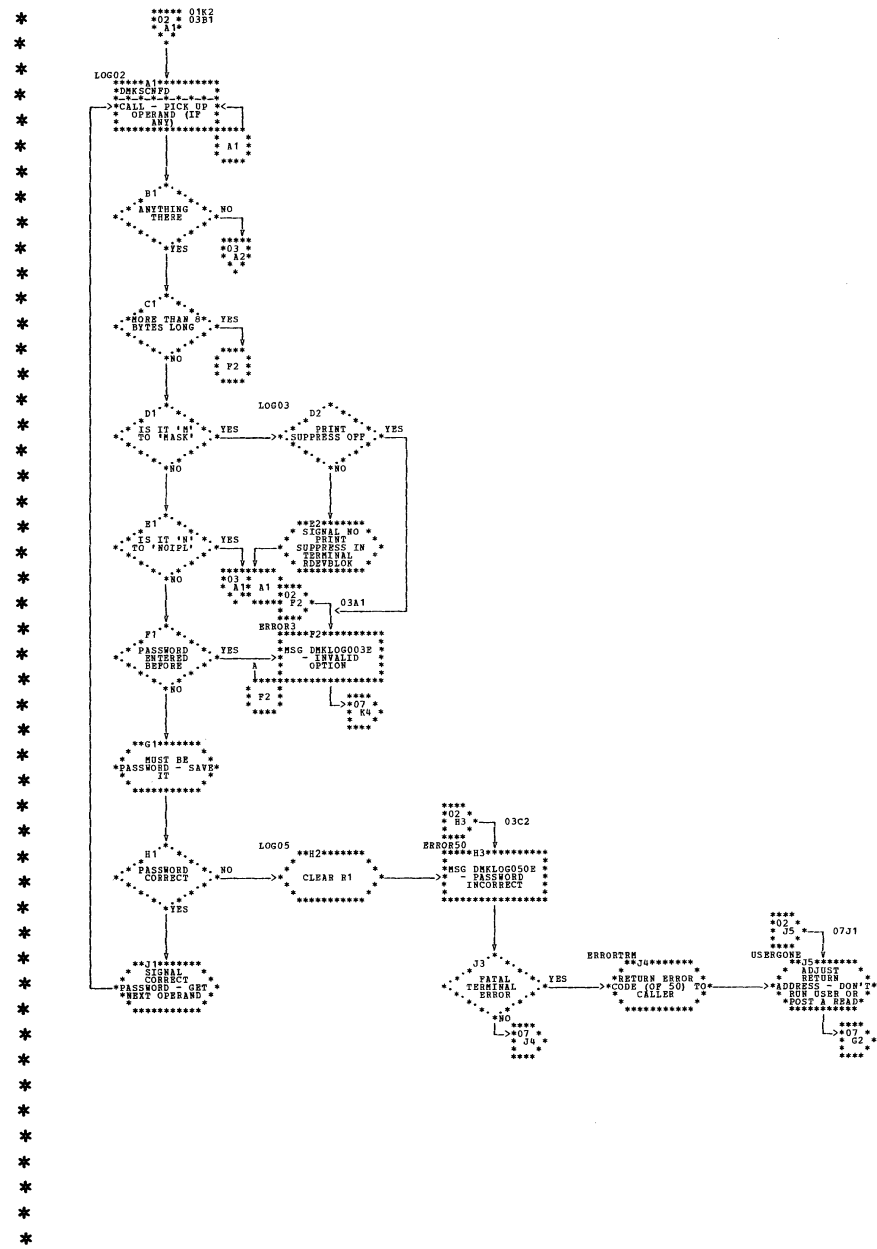
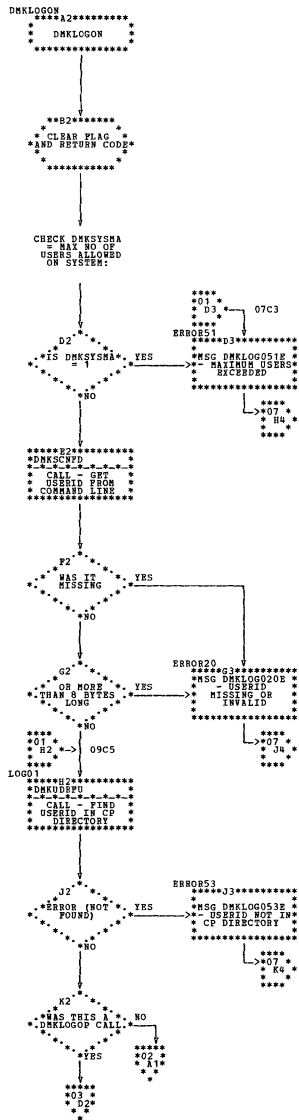






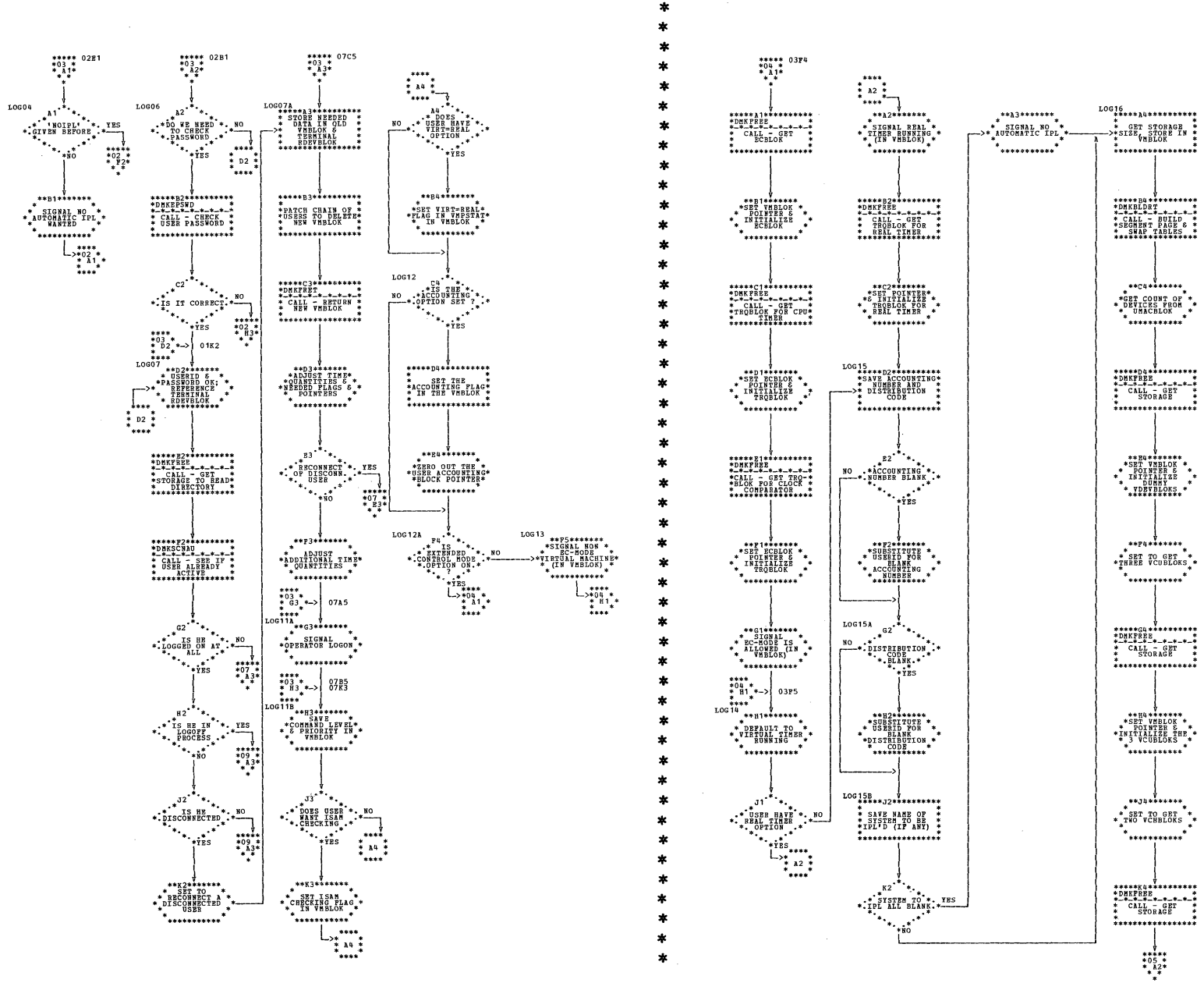


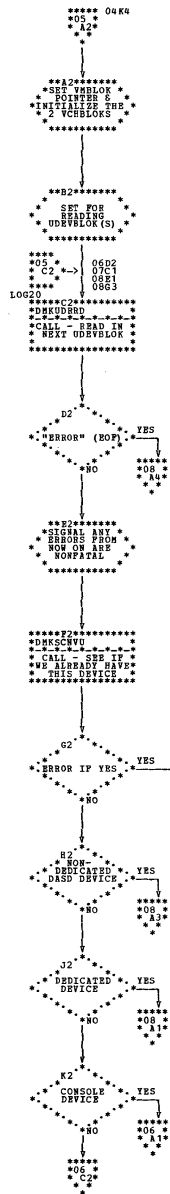




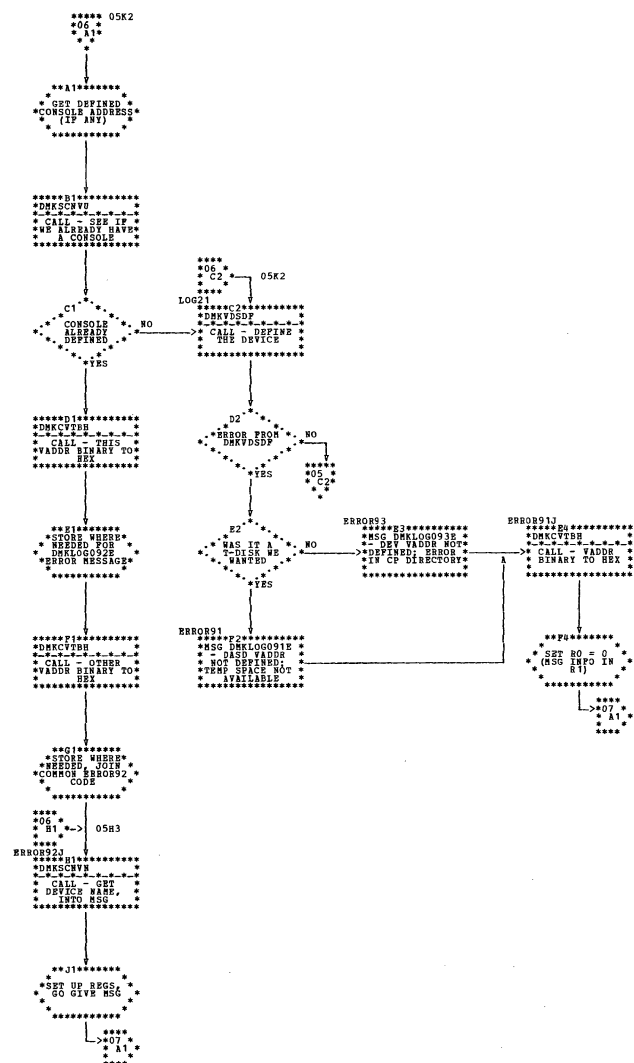
| DMKLOG -- Process LOGON/LOGIN Command; Logon the User or Operator (Parts 1 and 2 of 9)

DMKLOG -- Process LOGON/LOGIN Command; Logon the User or Operator (Parts 3 and 4 of 9)



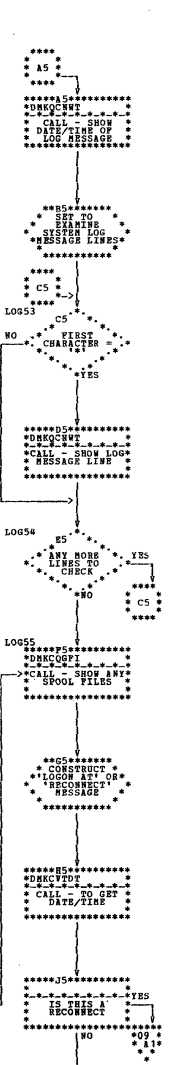
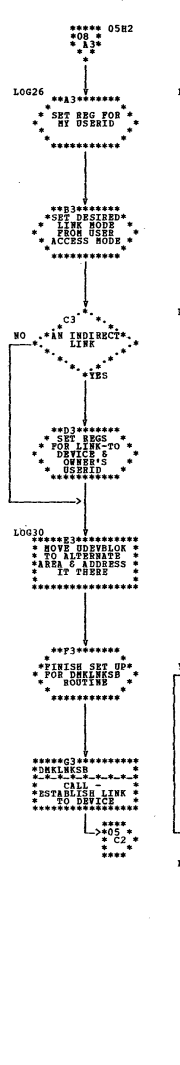
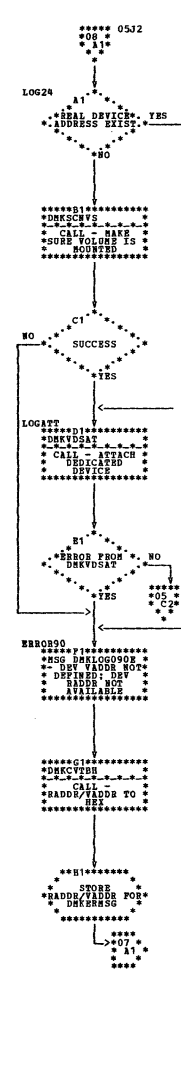
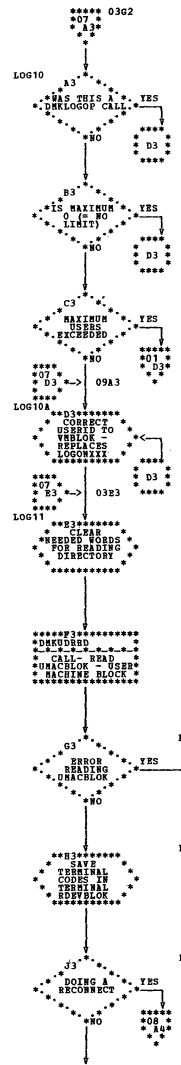
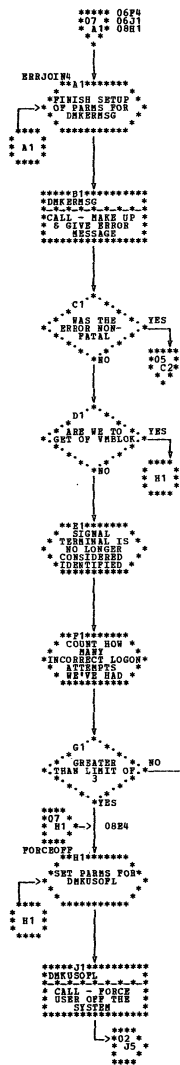


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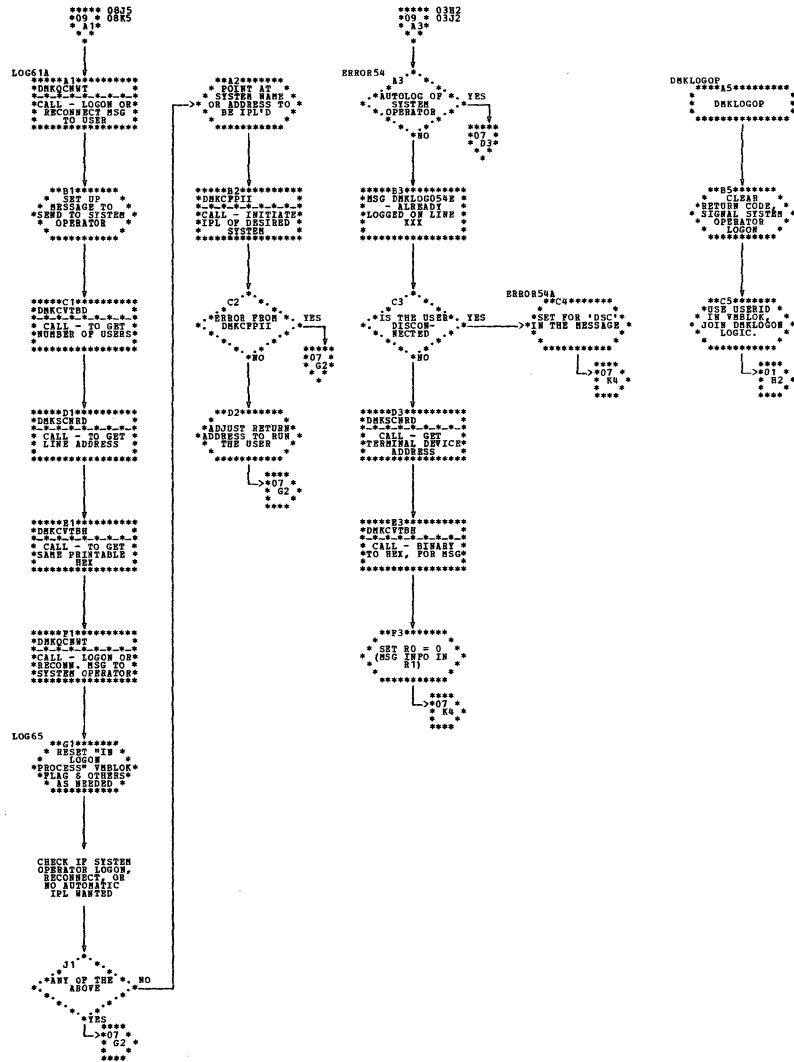


| DMKLOG -- Process LOGON/LOGIN Command; Logon the User or Operator (Parts 5 and 6 of 9)

DMKLOG -- Process LOGON/LOGIN Command; Logon the User or Operator (Part 7 and 8 of 9)

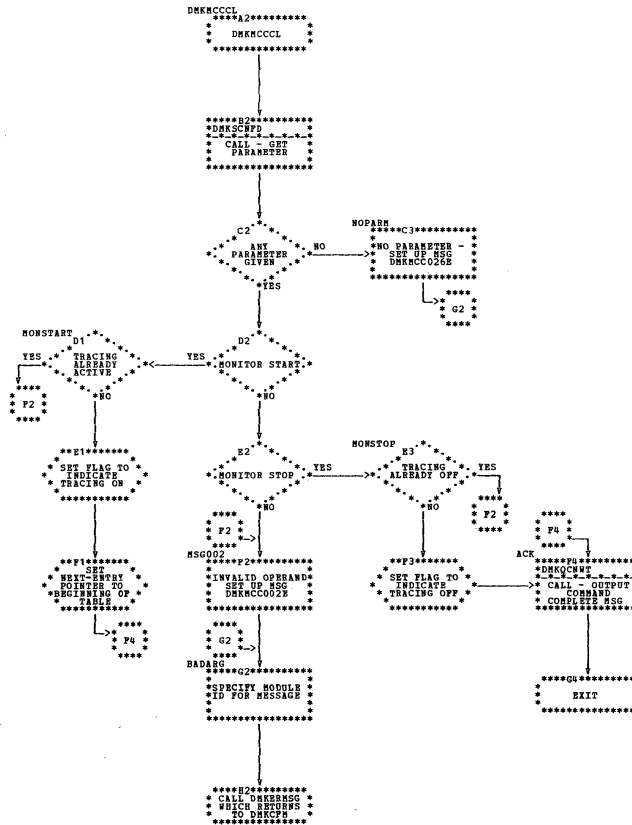


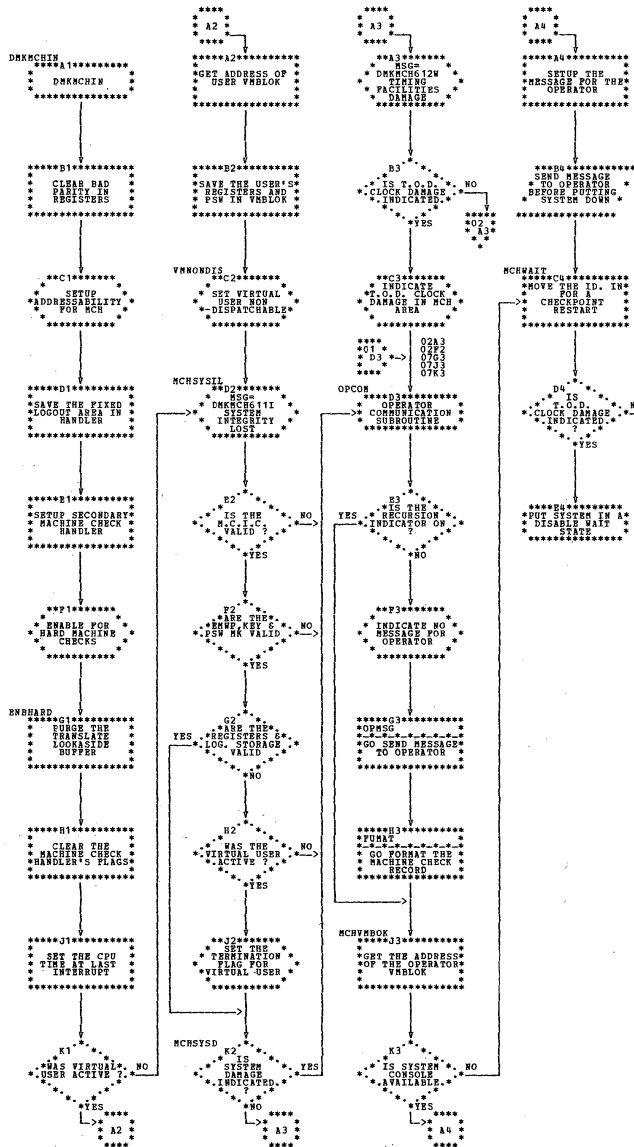
\* \* \* \* \*



| DMKLOG -- Process LOGON/LOGIN Command; Logon the User or Operator (Part 9 of 9)

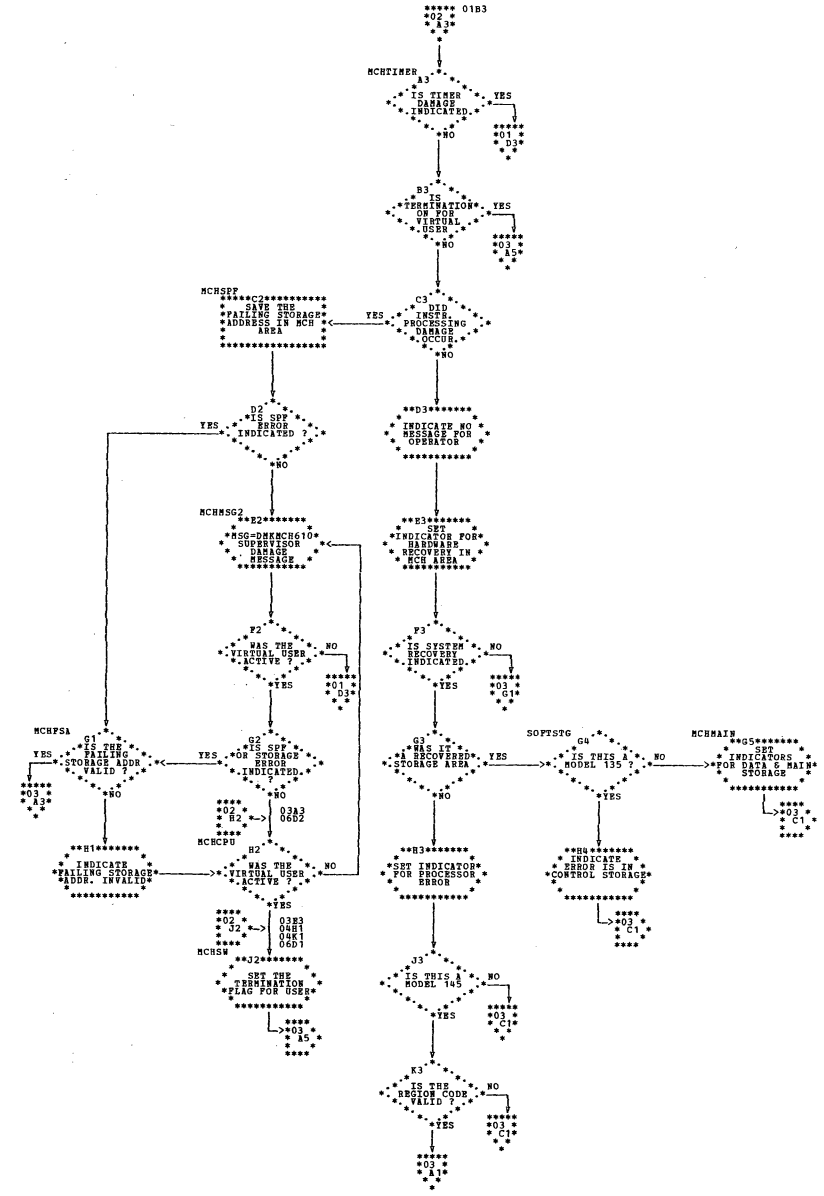
| DMKMCC -- VM Monitor Command Handler (Part 1 of 1)



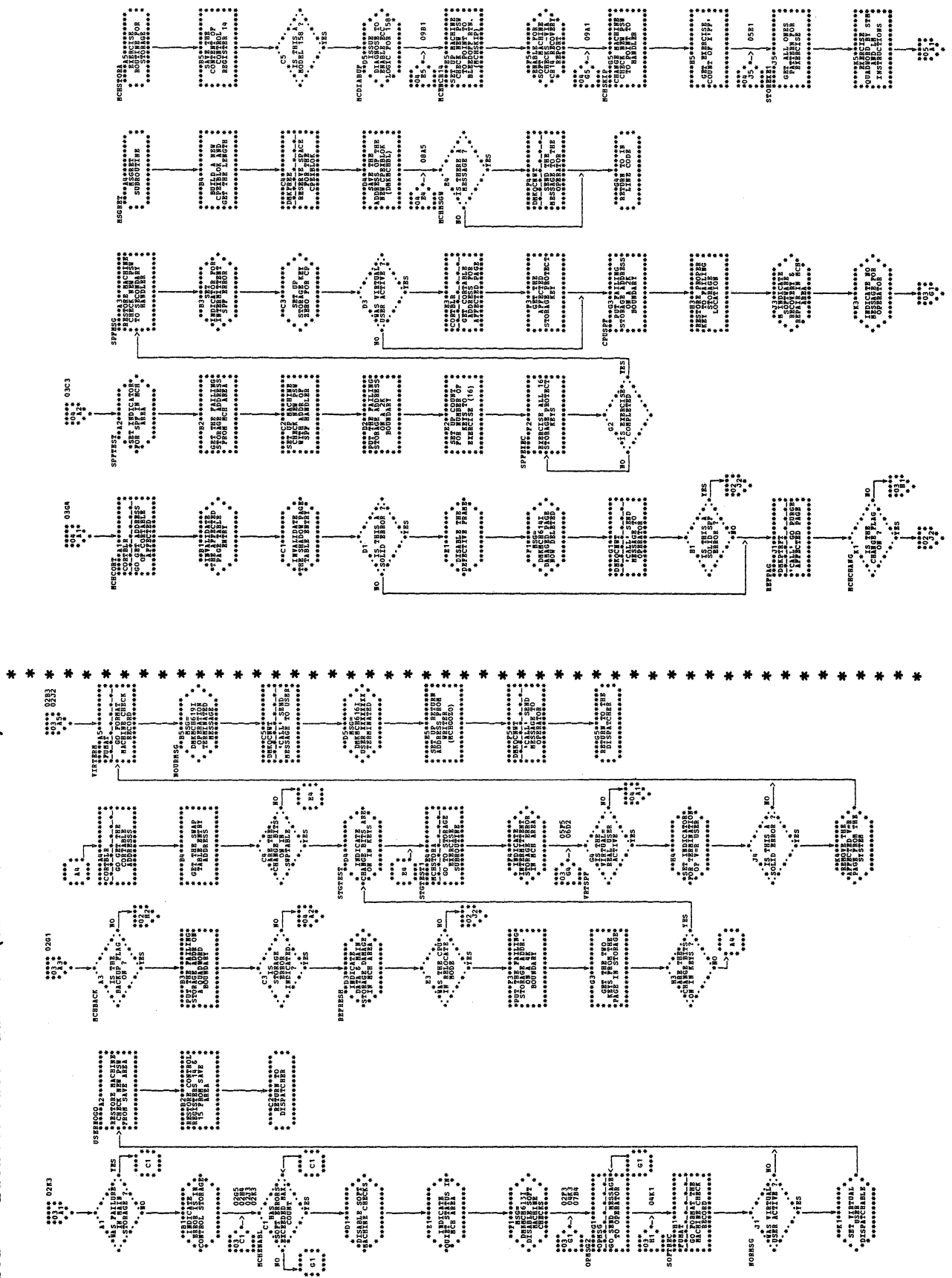


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DMKMCH -- Machine Check Handler (Parts 1 and 2 of 9)

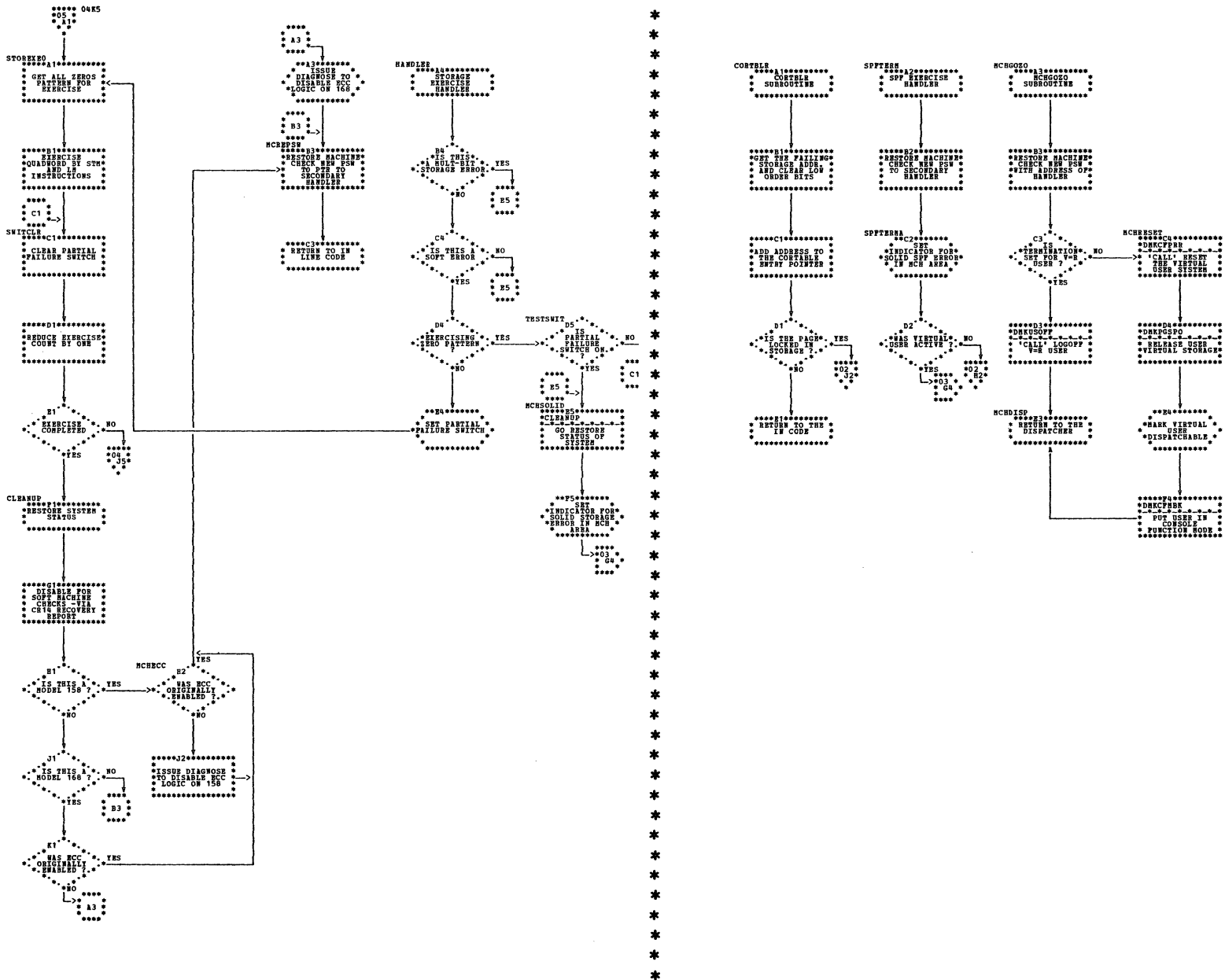


DKMCH -- Machine Check Handler (Parts 3 and 4 of 9)



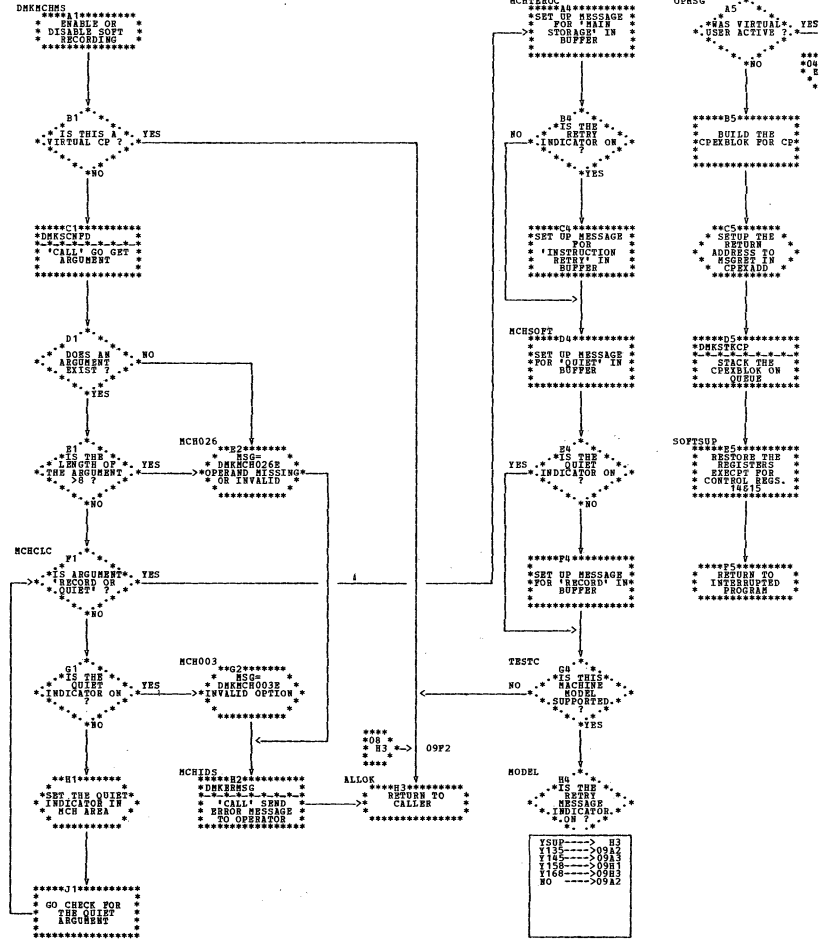
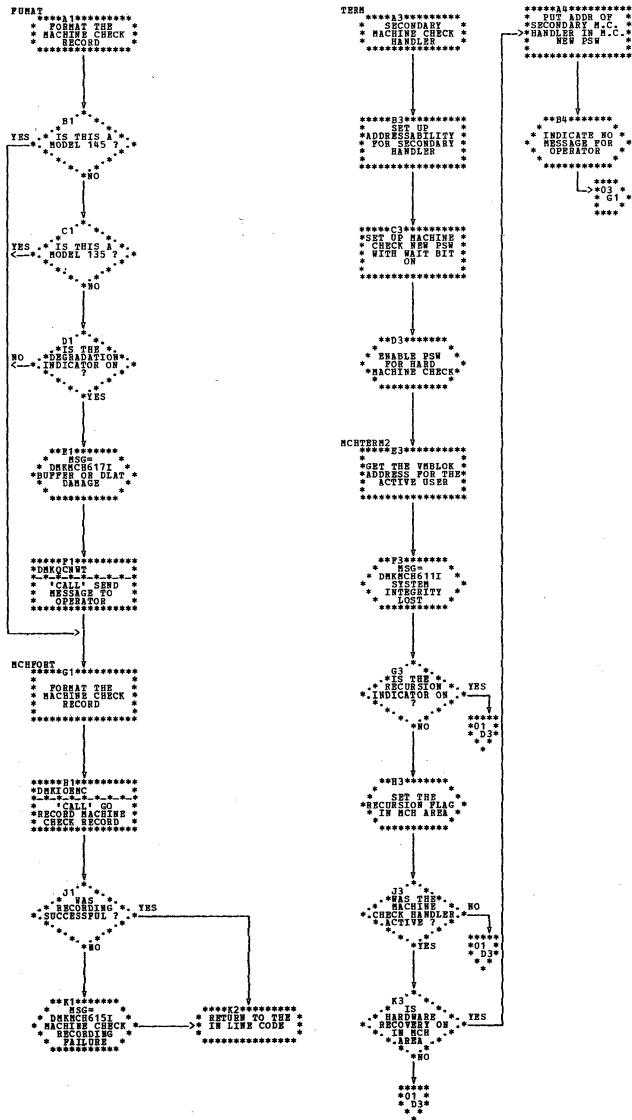
\* \* \* \* \*





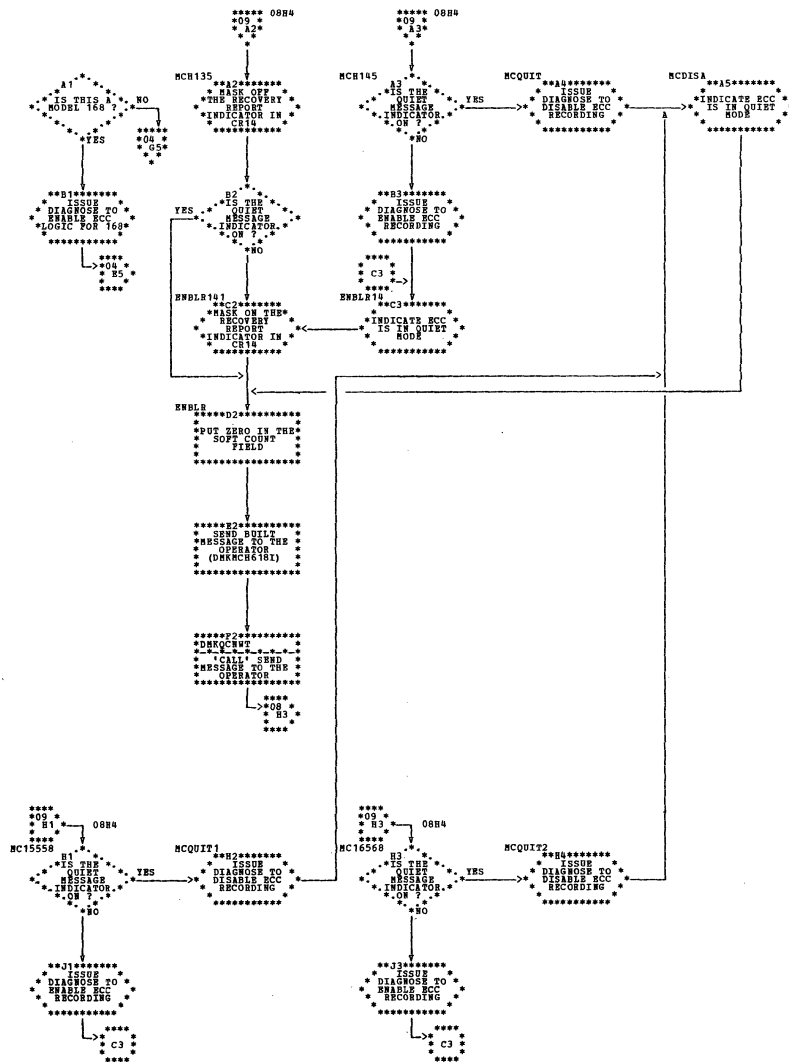
| DMKMCH -- Machine Check Handler (Parts 5 and 6 of 9)

| DMKMCH -- Machine Check Handler (Parts 7 and 8 of 9)

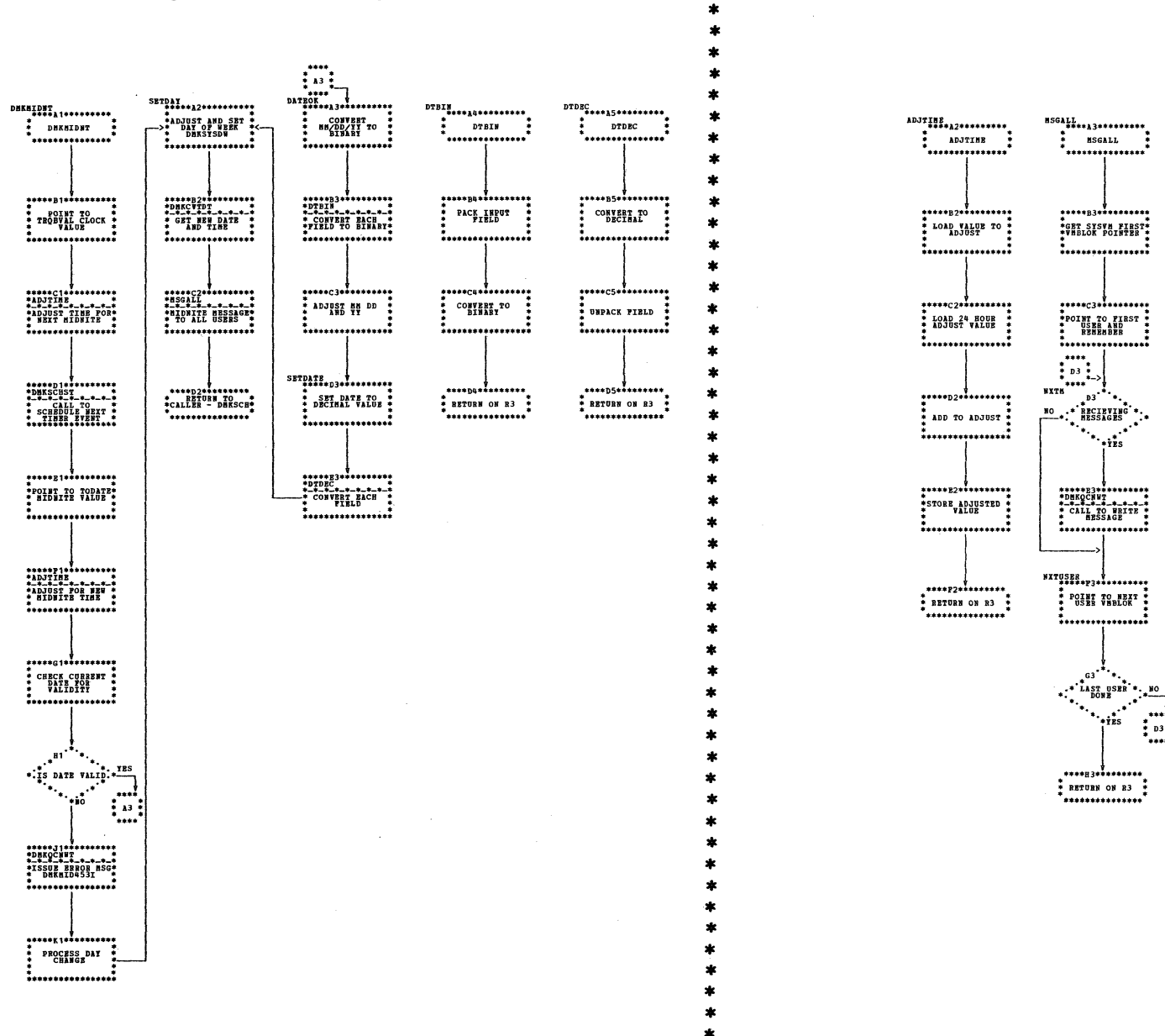


YSHR	---	H3
Y158	---	09A2
Y159	---	09B1
Y160	---	09B1
Y161	---	09B1
Y162	---	09A2
NO	---	09A2

\*\*\*\*\*

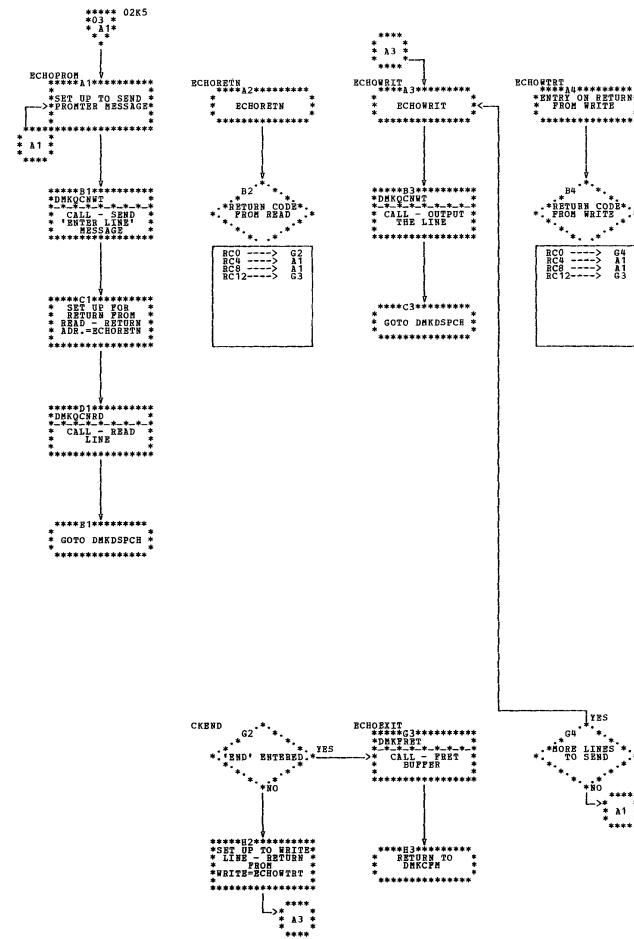


DMKID -- Change Date at Midnight (Parts 1 and 2 of 2)

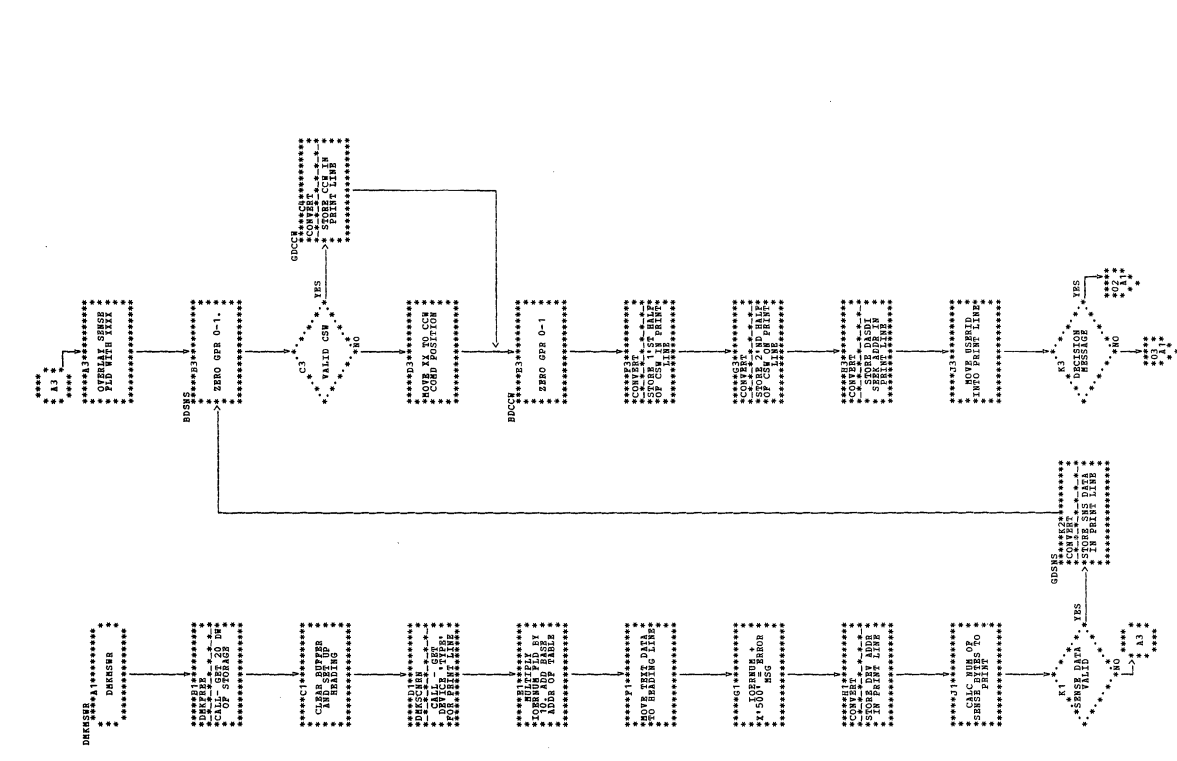
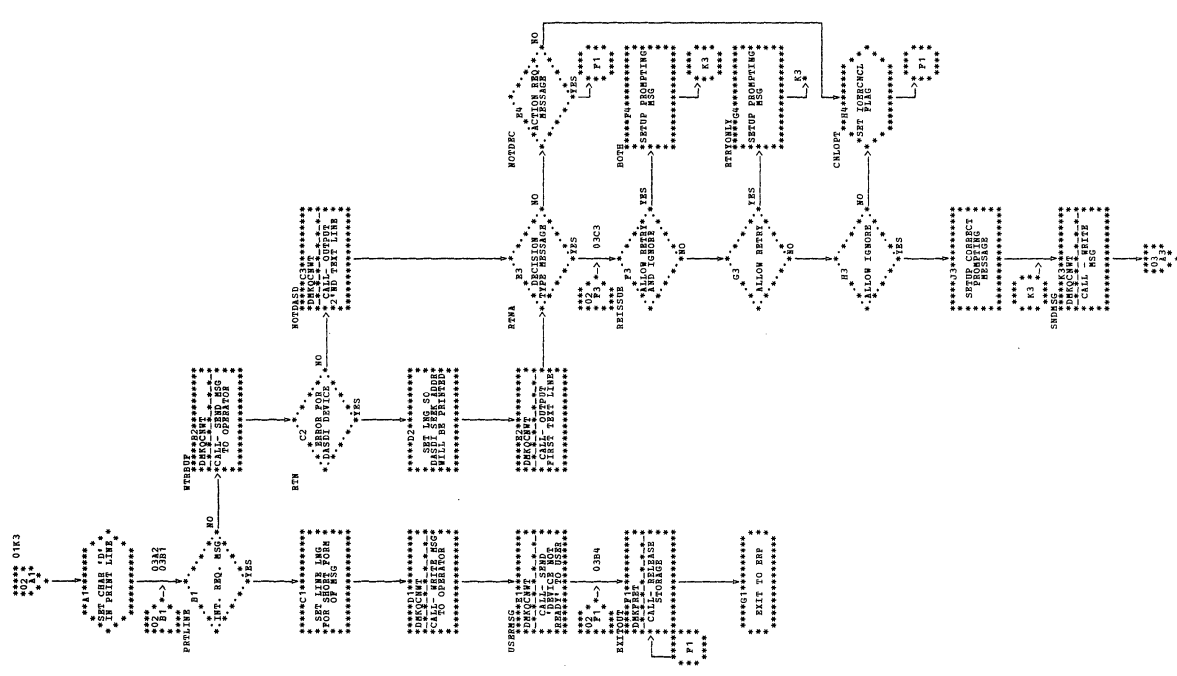




DMKMSG -- Process ECHO, MSG, and WNG Commands (Part 3 of 3)

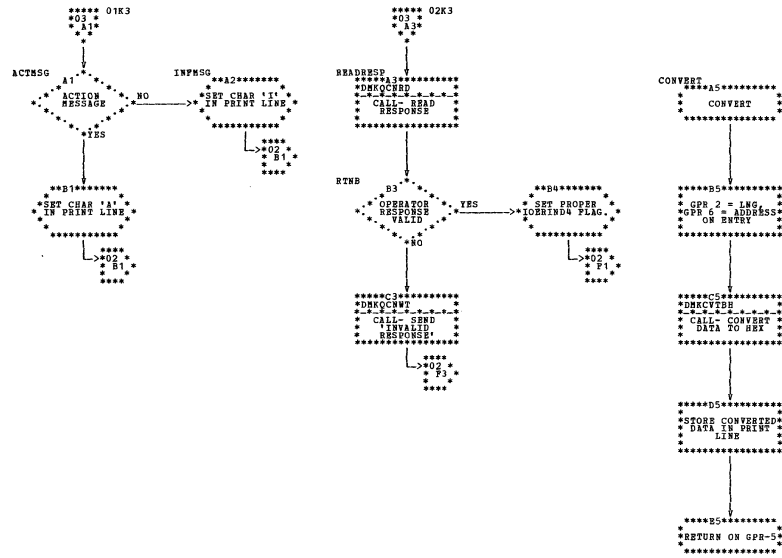


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DKMSW -- ERP Message Writer (Parts 1 and 2 of 3)

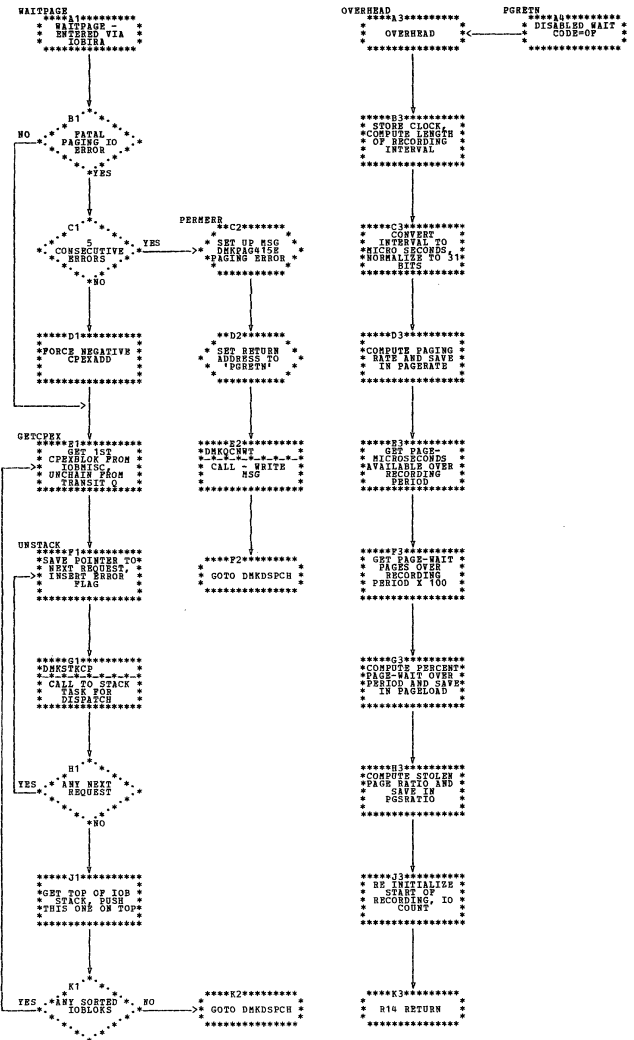
DMKMSW -- ERP Message Writer (Part 3 of 3)



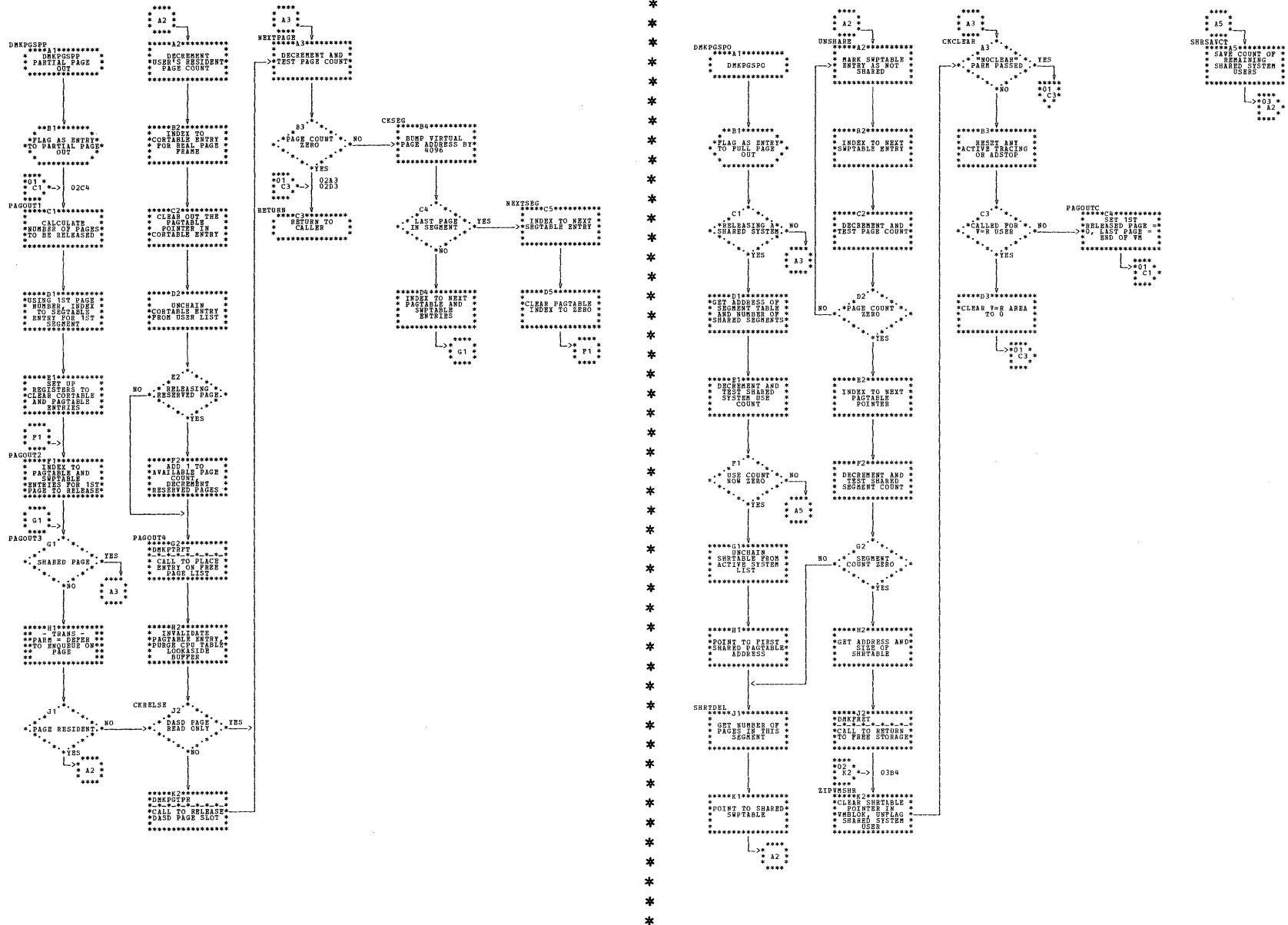


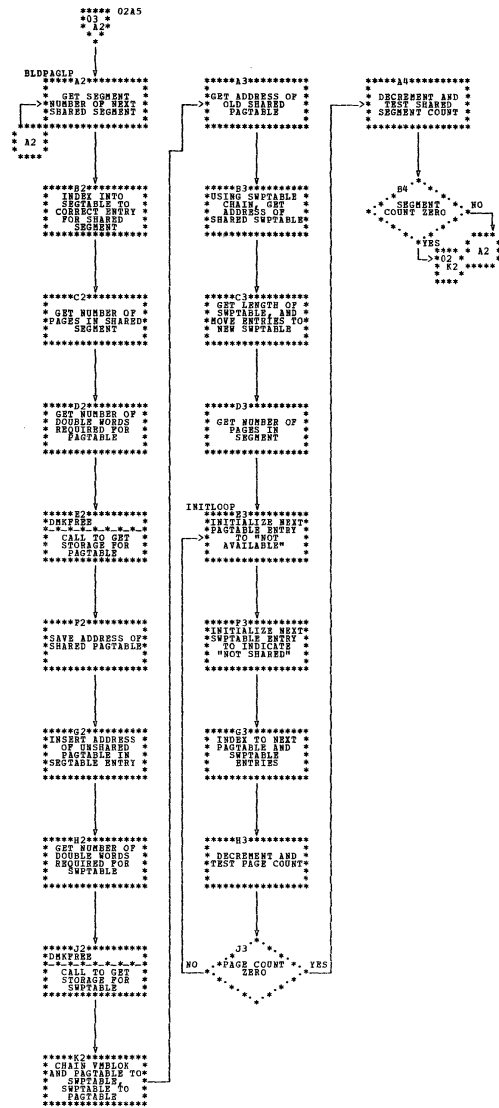




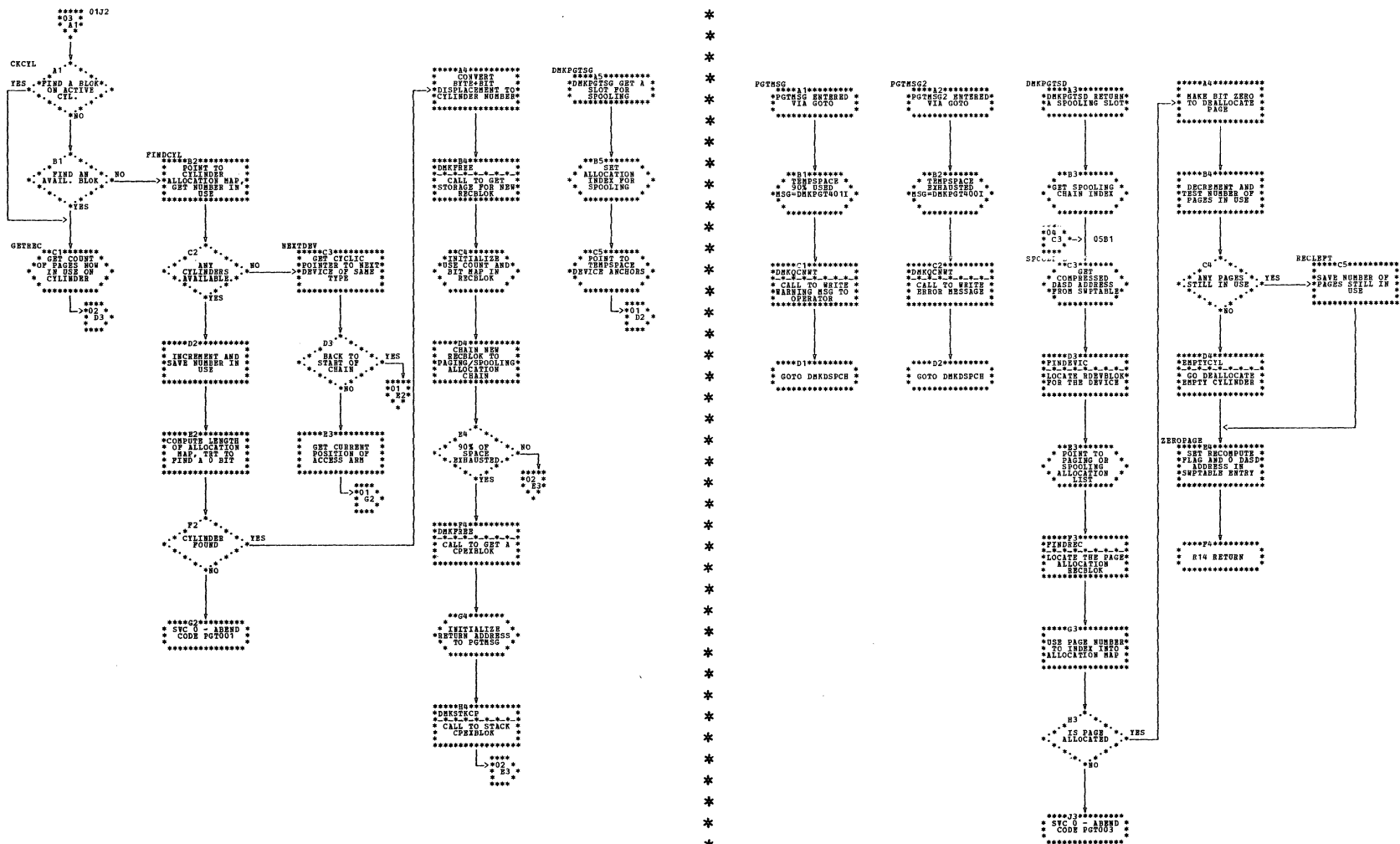


DMKPGS -- Release Virtual Storage (Parts 1 and 2 of 3)



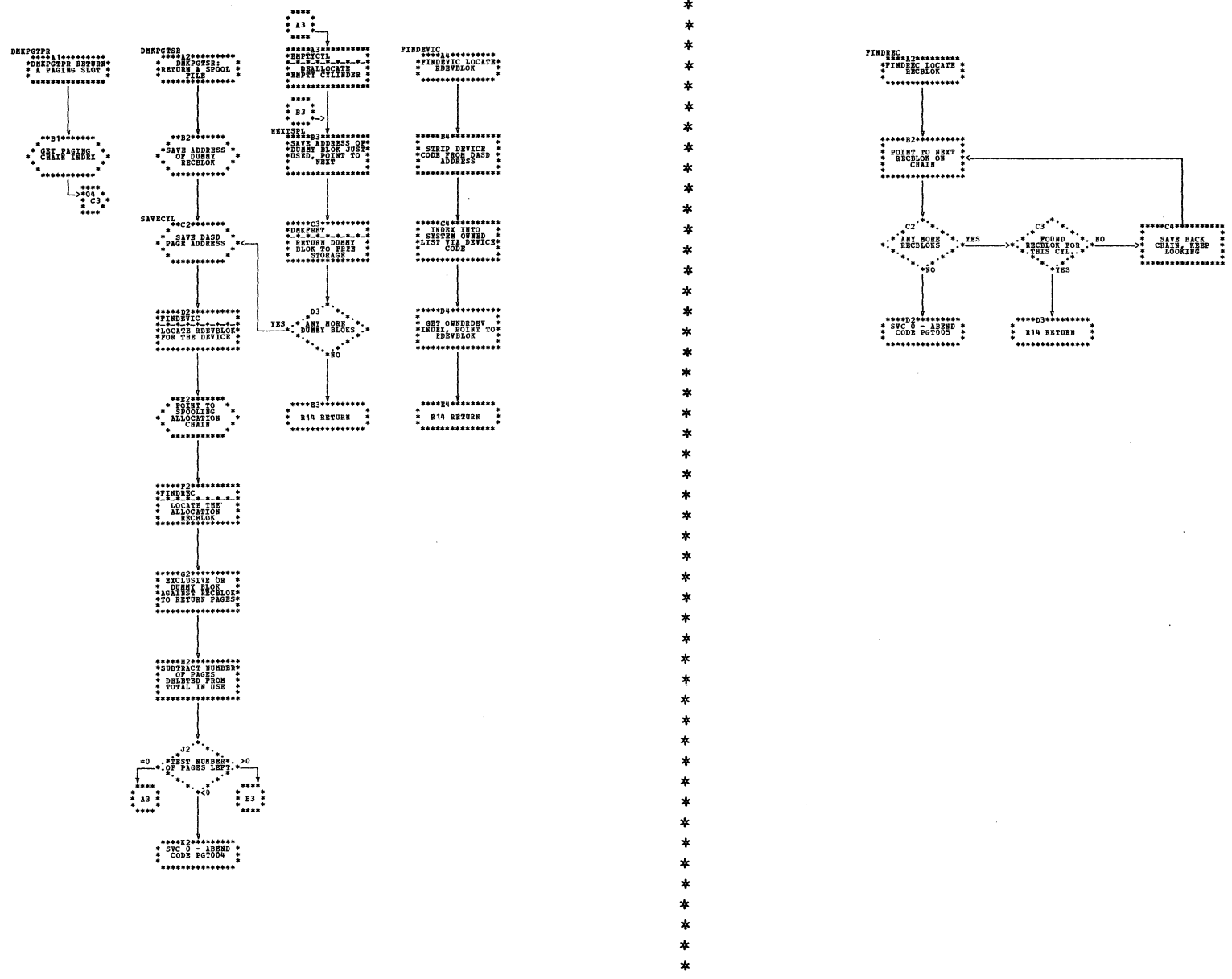






DMKPGT -- Allocate DASE/Virtual Storage (Parts 3 and 4 of 8)

DMKPGT -- Allccate DASD/Virtual Storage (Parts 5 and 6 of 8)





DNKPCTVL  
\*\*\*\*\*  
DNKPCTVL  
\*\*\*\*\*  
DNKPCTVL  
\*\*\*\*\*  
DNKPCTVL  
\*\*\*\*\*

\*\*\*\*\*B1\*\*\*\*\*  
\*GET POINTER TO  
\* CYLINDER  
\* ALLOCATION BLOC\*

\*\*\*\*\*C\*\*\*\*\*  
\*CALCULATE  
\* DISPLACE  
\* DISPLACE  
\* FOR CYLINDER TO  
\* DEALLOCATE\*

\*\*\*\*\*D1\*\*\*\*\*  
\*POINT TO  
\* CORRECT BYTE  
\* TEST IF  
\* CYLINDER IS  
\* ALLOCATED\*

\*\*\*\*\*E1\*\*\*\*\*  
\*IS CYLINDER  
\* ALLOCATED\*

\*\*\*\*\*F1\*\*\*\*\*  
\*SVC 0 - ARRD  
\* CODE PG7006\*

\*\*\*\*\*E2\*\*\*\*\*  
\*DEFERRED  
\* PAGING VOLUME\*

\*\*\*\*\*F2\*\*\*\*\*  
\*LAST  
\* CYLINDER ON  
\* CHAIN\*

UNCHAIN  
\*\*\*\*\*F3\*\*\*\*\*  
\*SET BIT TO ZERO  
\* TO FLAG  
\* CYLINDER  
\* DEALLOCATED\*

\*\*\*\*\*G2\*\*\*\*\*  
\*DECREMENT  
\* CYLINDER USE  
\* COUNT FOR  
\* DEVICE AND  
\* SYSTEM\*

\*\*\*\*\*H3\*\*\*\*\*  
\*UNLINK THE  
\* BLOC FROM  
\* CHAIN FOR ALLOCATION\*

\*\*\*\*\*J3\*\*\*\*\*  
\*DECREMENT  
\* THE BLOC TO  
\* THE STORAGE\*

\*\*\*\*\*K3\*\*\*\*\*  
\*R14 RETURN\*

DNKPCTVY  
\*\*\*\*\*  
DNKPCTVY  
\*\*\*\*\*  
DNKPCTVY  
\*\*\*\*\*

\*\*\*\*\*M\*\*\*\*\*  
\*TEST ADDRESS  
\* FREED TO  
\* LOCATE IN  
\* UNUSED VIRTUAL  
\* PAGE\*

\*\*\*\*\*N\*\*\*\*\*  
\*IS PAGES  
\* AVAILABLE\*

\*\*\*\*\*O\*\*\*\*\*  
\*SET BIT NO. 1 TO  
\* ALLOCATE PAGE\*

\*\*\*\*\*P\*\*\*\*\*  
\*ADD BYTE \* BIT  
\* TO BLOC TO GET  
\* RELATIVE  
\* PAGE NUMBER\*

\*\*\*\*\*Q\*\*\*\*\*  
\*MULTIPLY NUMBER  
\* BY 8096 TO GET  
\* PHYSICAL ADDRESS\*

\*\*\*\*\*R\*\*\*\*\*  
\*SET VIRTUAL  
\* ADDRESS INTO  
\* CALLER'S GPR\*

\*\*\*\*\*S\*\*\*\*\*  
\*GPR  
\* 0883  
\* R4 ->  
\* \*\*\*\*\*  
\* R14 RETURN  
\* \*\*\*\*\*

\*\*\*\*\*T\*\*\*\*\*  
\*EXIT  
\* \*\*\*\*\*  
\* R14 RETURN  
\* \*\*\*\*\*

\*\*\*\*\*U\*\*\*\*\*  
\*R14 RETURN\*

ORREQUEST  
\*\*\*\*\*V\*\*\*\*\*  
\*CALL TO GET A  
\* CPEXBLOC\*

\*\*\*\*\*W\*\*\*\*\*  
\*SAVE BLOC AND  
\* R14 AS RETURN  
\* ADDRESS\*

\*\*\*\*\*X\*\*\*\*\*  
\*CHAIN CPEXBLOC  
\* TO END OF LIST  
\* OF DEFERRED  
\* REQUESTS\*

\*\*\*\*\*Y\*\*\*\*\*  
\*GOTO DNKDSPCR\*

DNKPCTVZ  
\*\*\*\*\*  
DNKPCTVZ  
\*\*\*\*\*  
DNKPCTVZ  
\*\*\*\*\*

\*\*\*\*\*Z2\*\*\*\*\*  
\*INDEX INTO  
\* FREED LIST  
\* VIRTUAL PAGE  
\* NUMBER\*

\*\*\*\*\*Z4\*\*\*\*\*  
\*TEST BIT IN  
\* BYTE TO SEE IF  
\* PAGE IS  
\* ALLOCATED\*

\*\*\*\*\*Z5\*\*\*\*\*  
\*IS PAGE  
\* ALLOCATED\*

\*\*\*\*\*Z6\*\*\*\*\*  
\*SVC 0 - ARRD  
\* CODE PG7007\*

\*\*\*\*\*Z3\*\*\*\*\*  
\*GET POINTER TO  
\* LIST OF  
\* DEFERRED  
\* BLOC'S\*

\*\*\*\*\*Z7\*\*\*\*\*  
\*ANY ONE ON  
\* LIST\*

\*\*\*\*\*Z8\*\*\*\*\*  
\*CLEAR ADDRESS OF  
\* RELEASED PAGE  
\* TO LIST CPEXBLOC  
\* ON LIST\*

\*\*\*\*\*Z9\*\*\*\*\*  
\*DNKDPVZ  
\* \*\*\*\*\*  
\* CALL TO STACK  
\* CPEXBLOC  
\* \*\*\*\*\*

\*\*\*\*\*Z10\*\*\*\*\*  
\*R14 RETURN  
\* \*\*\*\*\*

RELEASE  
\*\*\*\*\*Z11\*\*\*\*\*  
\*SET BIT TO 0 TO  
\* DEALLOCATE THE  
\* PAGE\*

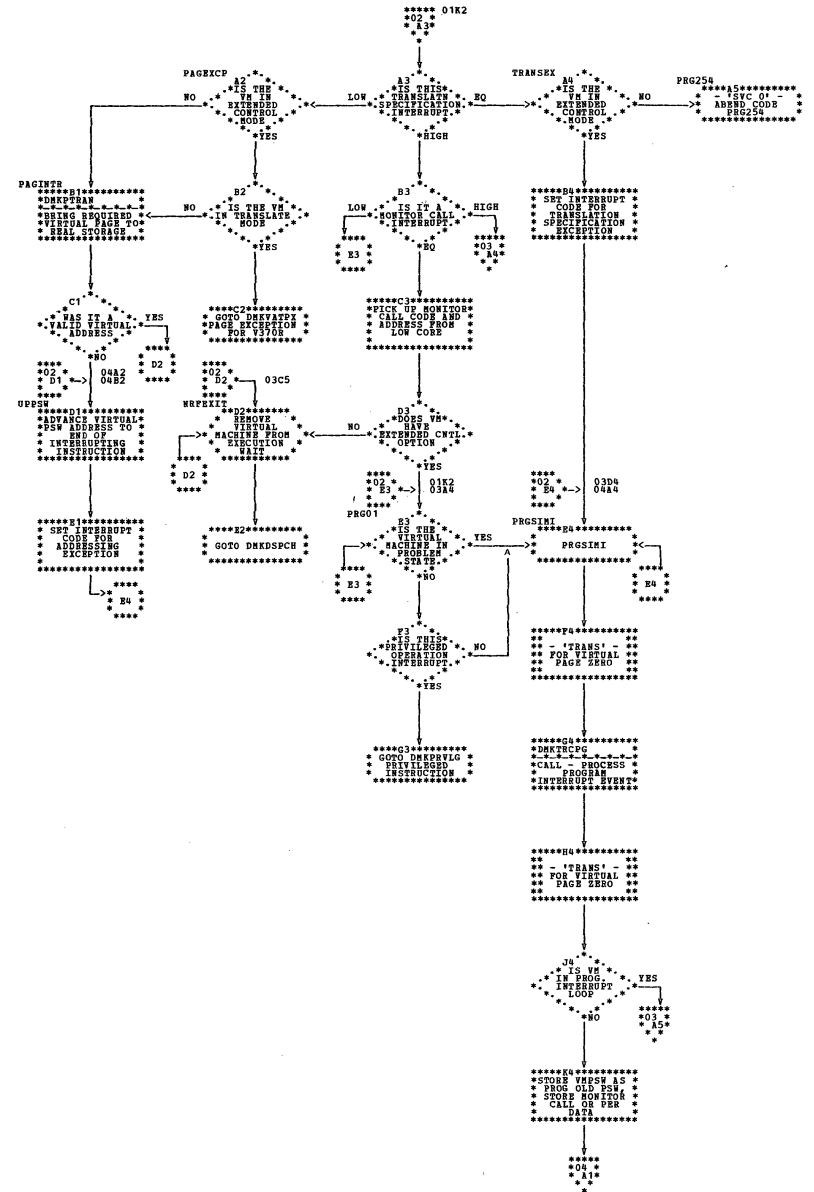
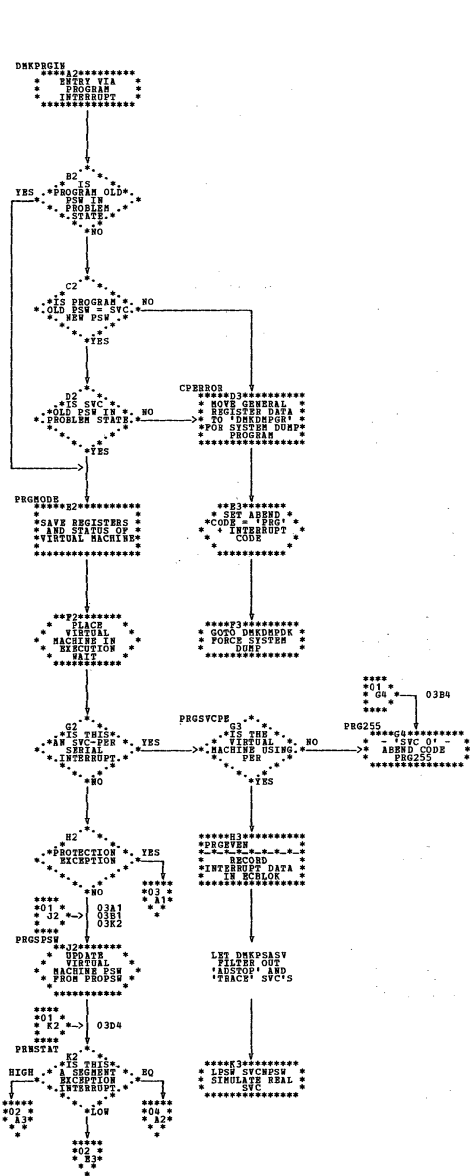
\*\*\*\*\*Z12\*\*\*\*\*  
\*R14 RETURN  
\* \*\*\*\*\*

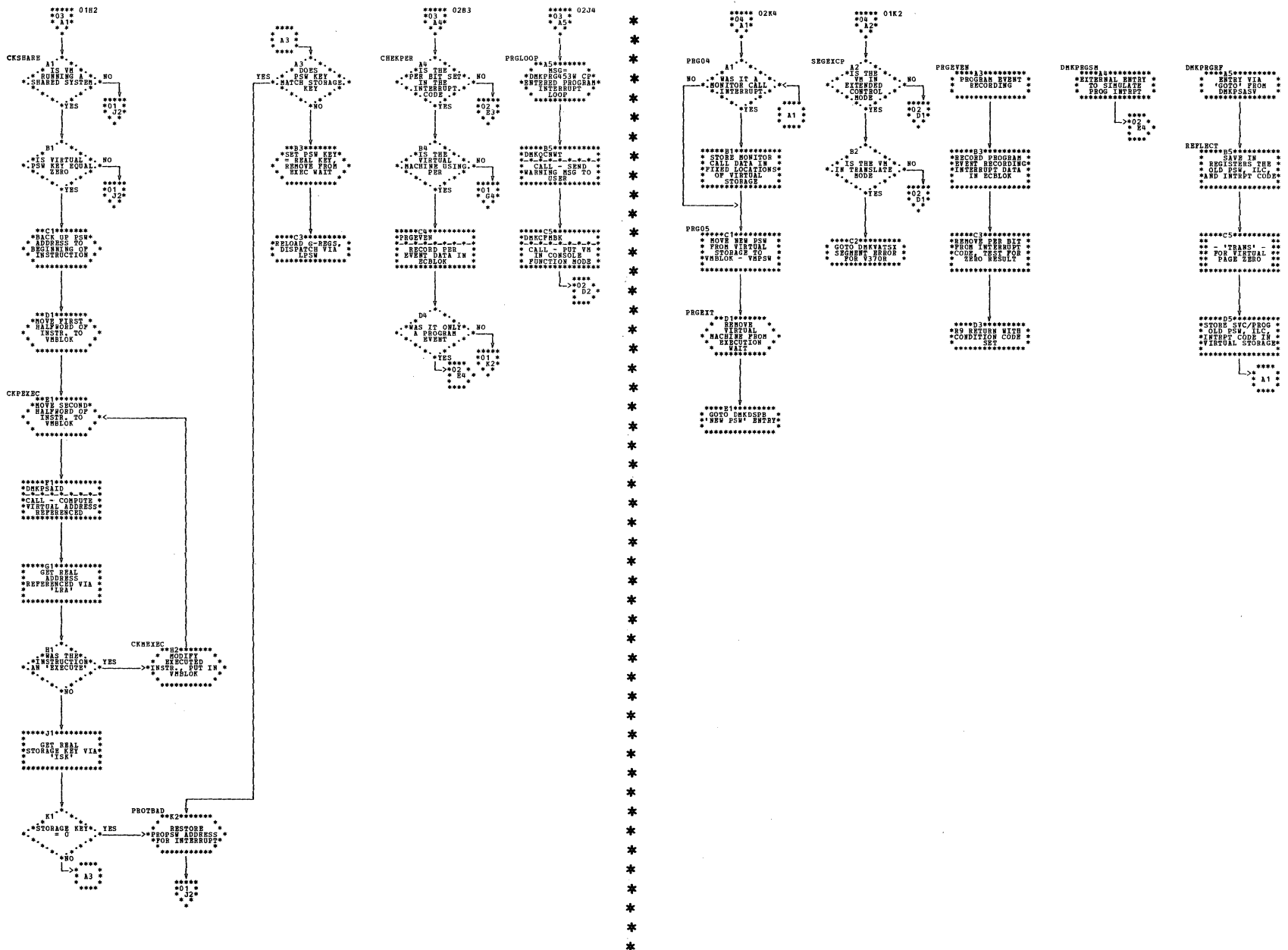
DNKPGT -- Allocate DASD/Virtual Storage (Parts 7 and 8 of 8)

SY20-0880-1, Page Modified by TNL SN20-2624, August 15, 1973

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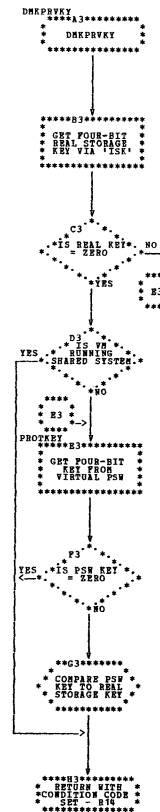
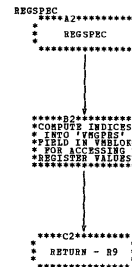
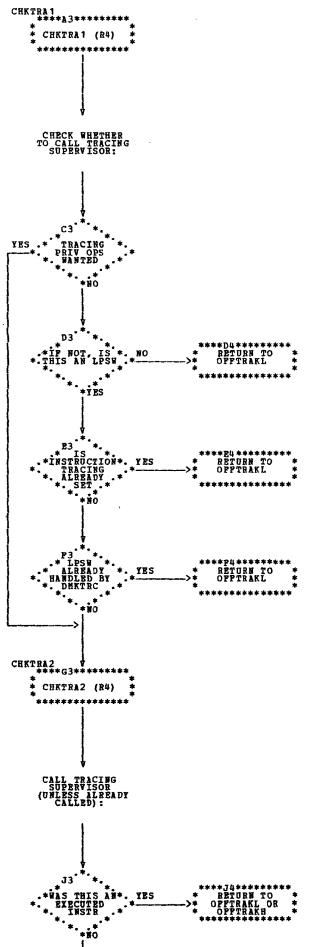
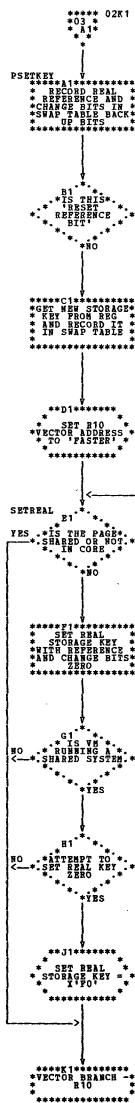
| DMKPRG -- Program Interrupt Handler (Parts 1 and 2 of 4)





DMKPRG -- Program Interrupt Handler (Parts 3 and 4 of 4)

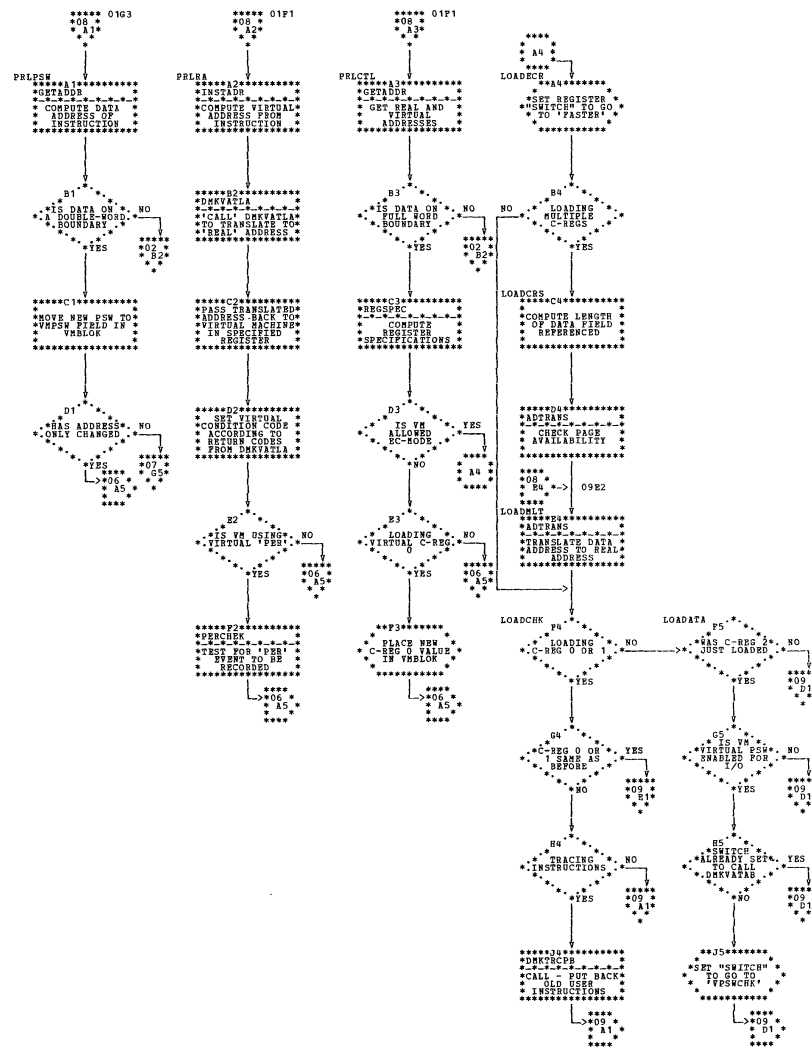
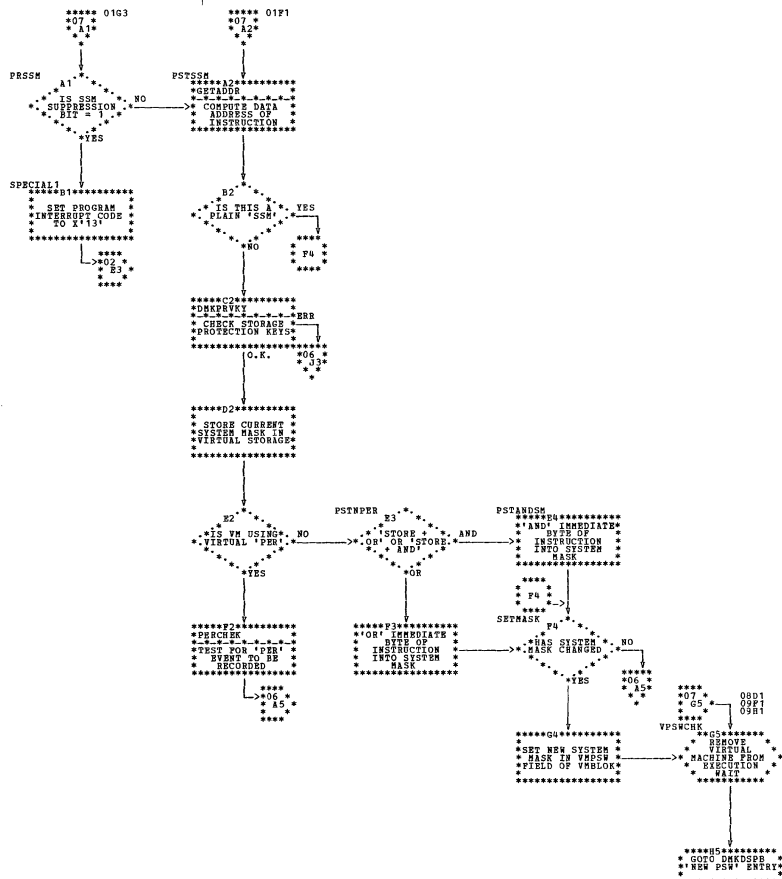




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| DMKPRV -- Simulate Virtual Machine Privileged Instruction (Parts 3 and 4 of 10)

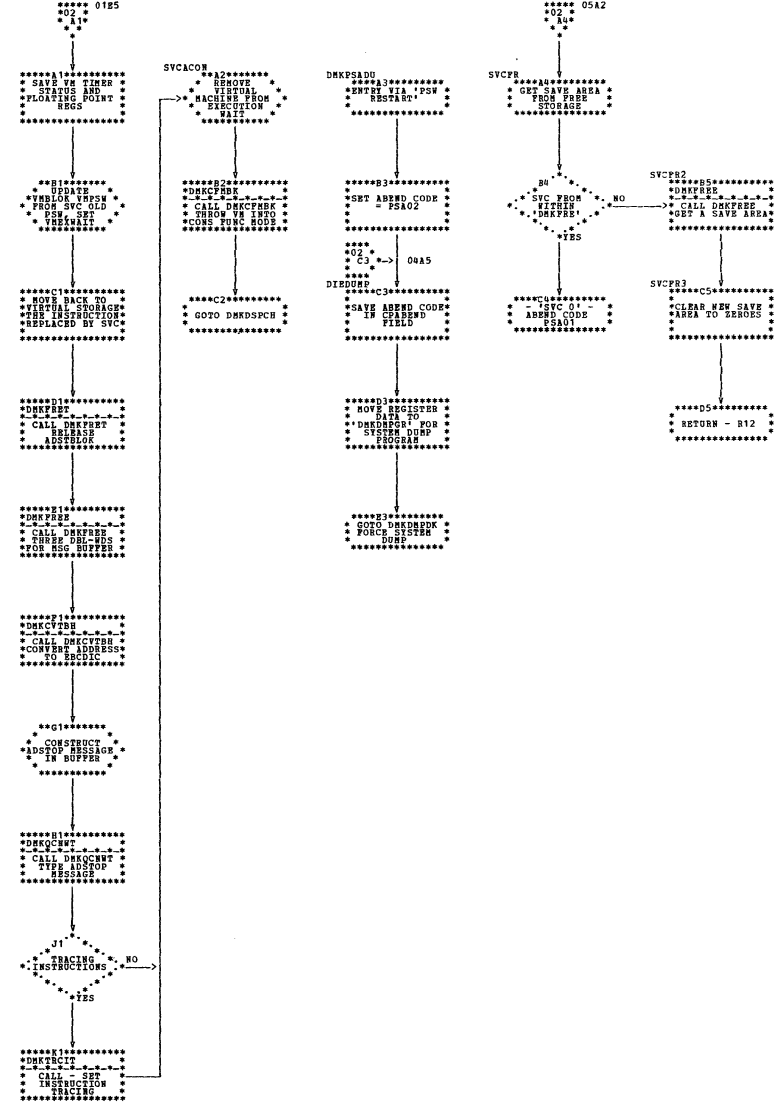
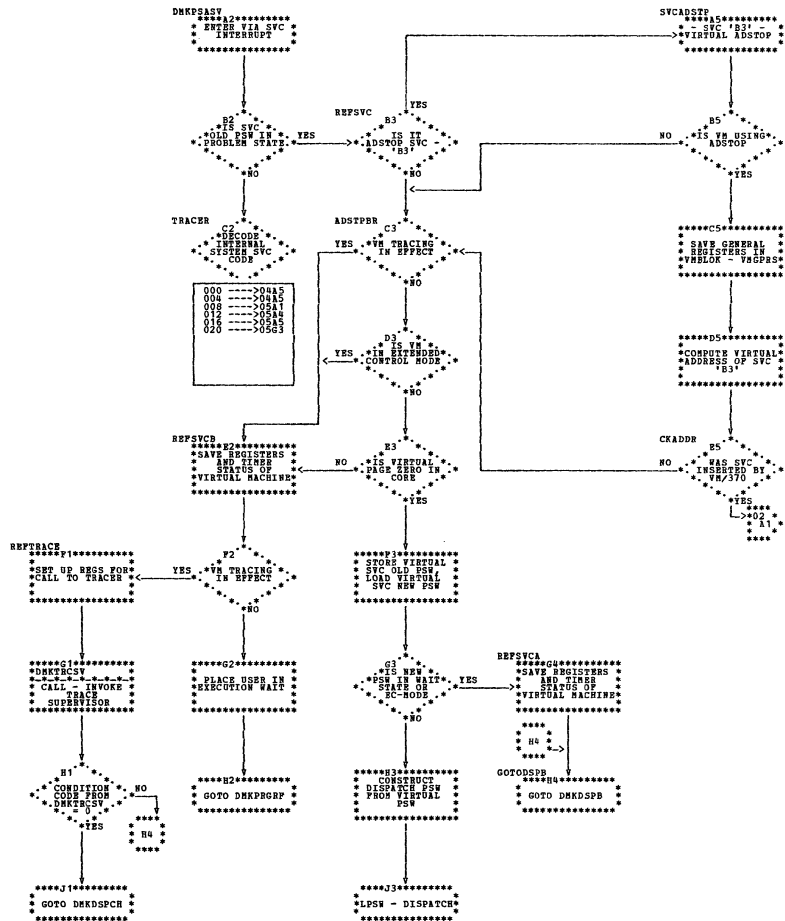




| DMKPRV -- Simulate Virtual Machine Privileged Instruction (Parts 7 and 8 of 10)

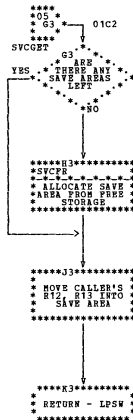
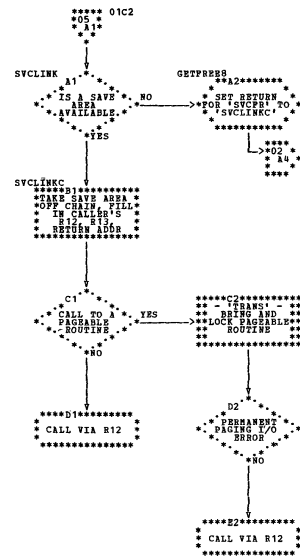




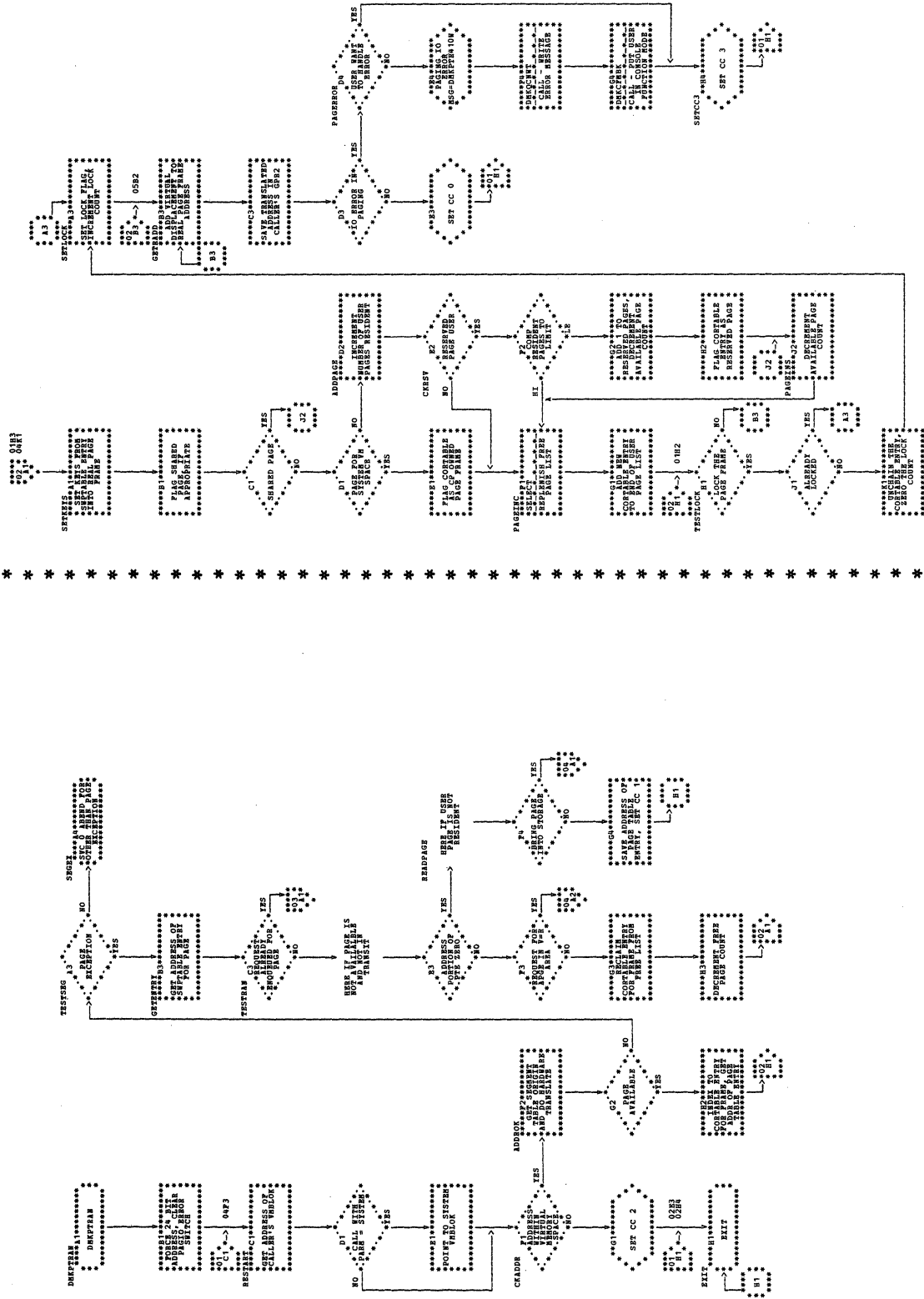


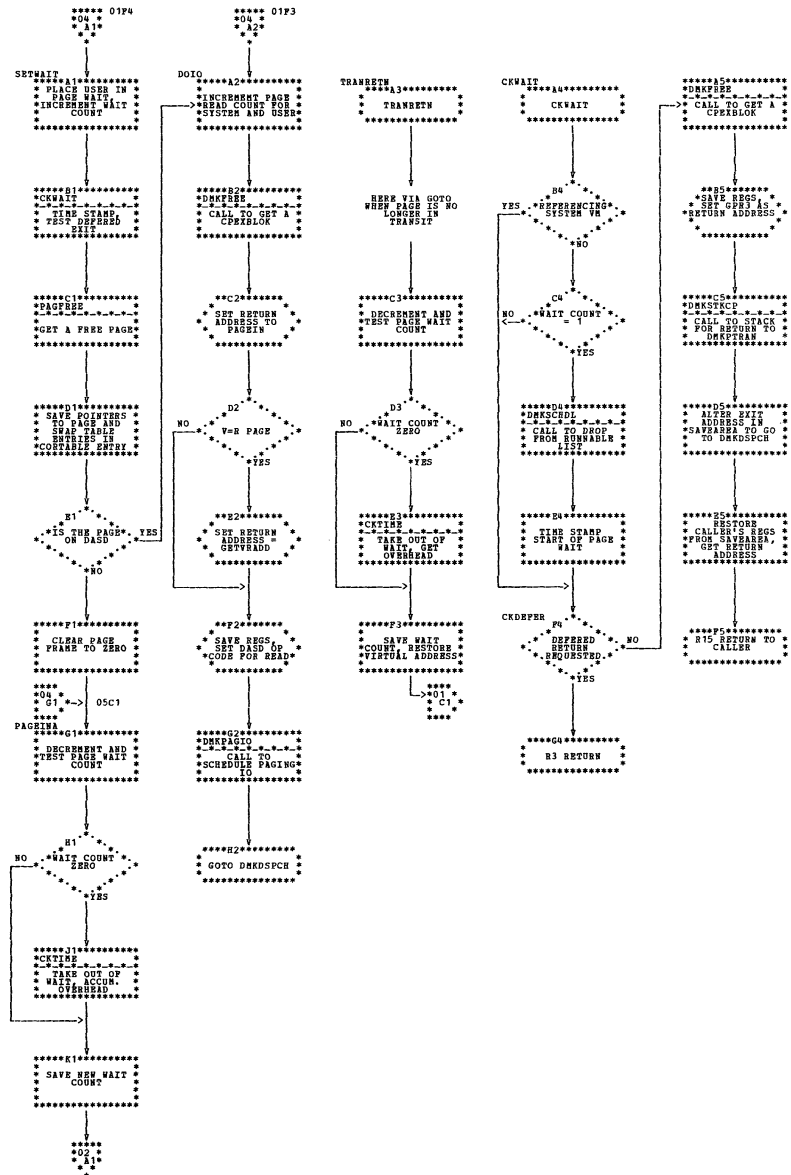
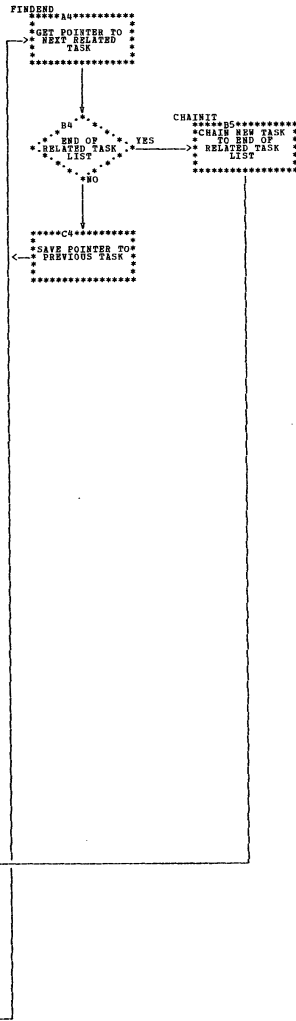
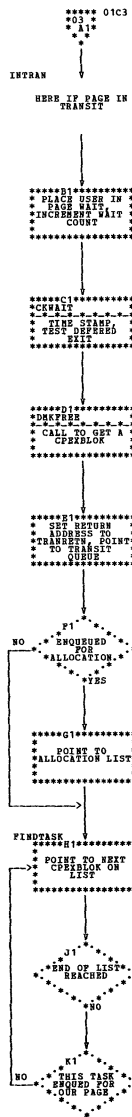
DMKPSA -- SVC and External Interrupt Handler (Parts 1 and 2 of 5)



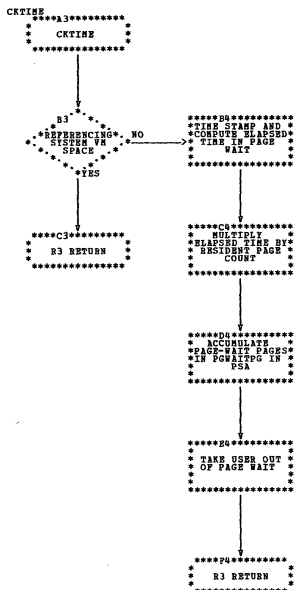
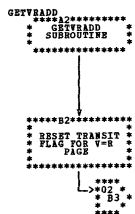


DMKPTR -- Real Storage Page Manager (parts 1 and 2 of 13)

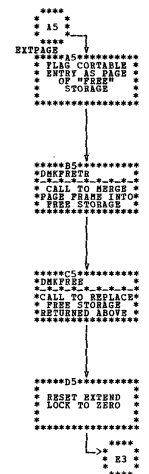
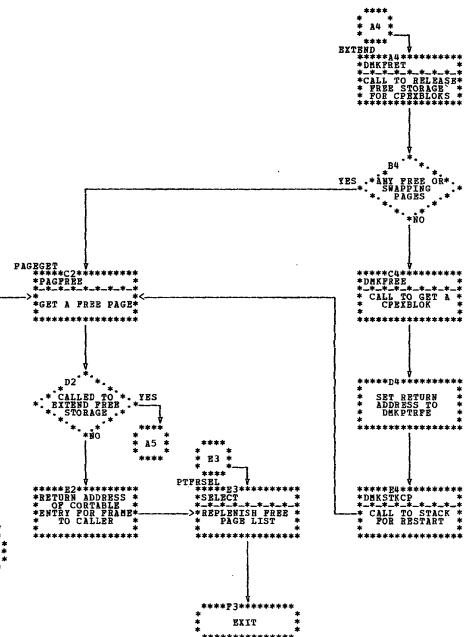
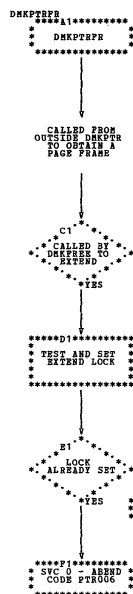


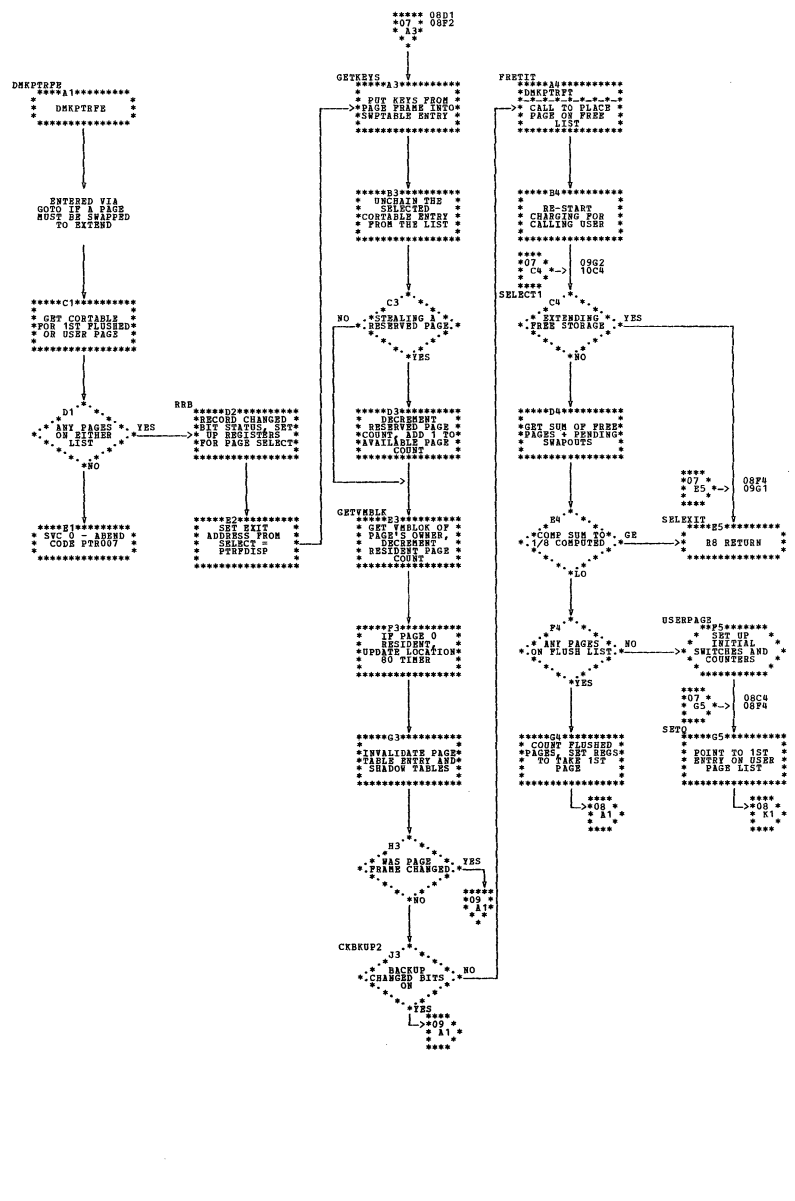


DMKPTR -- Real Storage Page Manager (Parts 5 and 6 of 13)

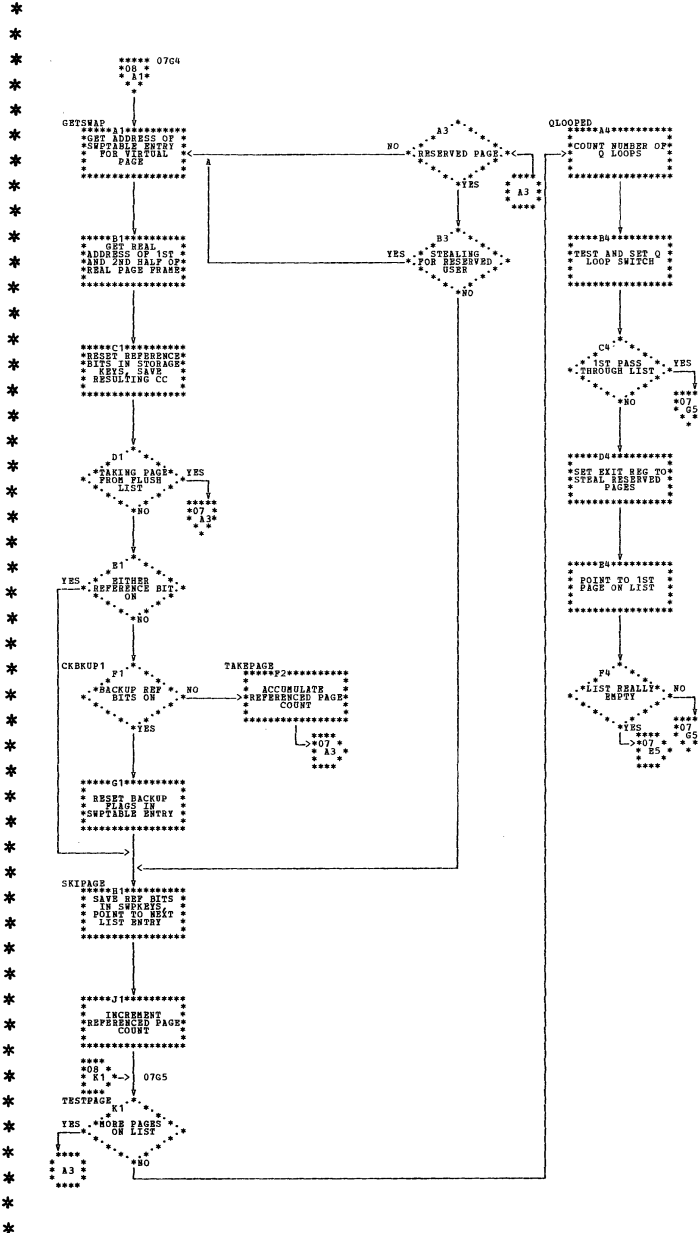


\* \* \* \* \*





DMKPTR -- Real Storage Page Manager (Parts 7 and 8 of 13)



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```

DMKPTRFD
*****A1*****
DMKPTRFD

```

```

ENTERED VIA
GOTO AFTER PAGE
WRITE COMPLETE

```

```

DECREMENT
NUMBER OF
PENDING
SWAPOUTS

```

```

D1
IS COMPLETE
OR
NO
YES
DOPRST
*****D2*****
DMKPTRFD
CALL TO PLACE
PAGE ON FREE
LIST

```

```

REWRITE
RESCHEDULE PAGE
FOR OUTPUT

```

```

SELDISP
*****
GOTO DMKSPCH

```

```

REWRITE
*****A3*****
REWRITE

```

```

INCREMENT
USER'S RESIDENT
PAGE COUNT

```

```

CHAIN CONTABLE
ENTER TO FORM
LIST VALIDATE
PAGE TABLE
ENTRY

```

```

D3
SET ADDRESS
IS-COMPUTE FLAG
IN SWPTABLE
ENTRY

```

```

R14 RETURN

```

```

DMKPTRFD
*****A4*****
DMKPTRFD

```

```

INCREMENT FREE
PAGE COUNT

```

```

C4
ANY DEFERRED
FREE PAGE
REQUESTS
NO
CHAINPAG
*****C5*****
CHAIN CONTABLE
ENTER TO END OF
FREE PAGE LIST

```

```

D4
REQUEST
FROM DMKPRSE
YES

```

```

REMOVE
ALLOCATION
ENQUEUE FLAG
FROM CHAINPAG
ENTRY

```

```

SAVEADDR
*****F4*****
SAVE ADDRESS OF
CONTABLE ENTRY
IN CPEXBLK

```

```

G4
VALID
POINTER TO
OLD PTE
NO
YES

```

```

H4
CLEAR PAGE
FRAME INDEX TO
ZERO

```

```

J4
MAKE NEXT
REQUIRED
REQUEST 1ST ON
LIST

```

\*\*\*\*\*

```

*****1104*****
*****A14*****
PTSTAN
DMKPRSD
CALL TO STACK
CPEXBLK FOR
WITH TO CALLER
E1
ANY RELATED
REQUESTS
YES
NO
*****1105*****
*****C1*****
PREEXIT
R14 RETURN

```

```

DMKPTRUL
*****A2*****
DMKPTRUL

```

```

UNLOCK A REAL
PAGE FRAME

```

```

C2
INDEX TO
CURRENT
CONTABLE ENTRY
VIA PAGE FRAME
NUMBER

```

```

D2
IS ENTRY
LOCKED
YES

```

```

E2
SVC 0 ABEND

```

```

D3
DECREMENT AND
TEST LOCK COUNT

```

```

E3
LOCK COUNT
< 0
YES

```

```

F3
SVC 0 ABEND

```

```

E5
SAVE LOCK COUNT

```

```

F5
LOCK COUNT
> 0
YES

```

```

G5
RESET LOCK
FLAG CHAIN
ENTRY TO TOP OF
USER LIST

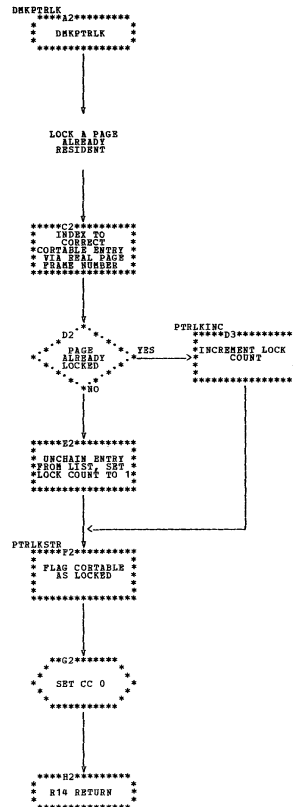
```

```

UNLOCKIT
*****H5*****
R14 RETURN

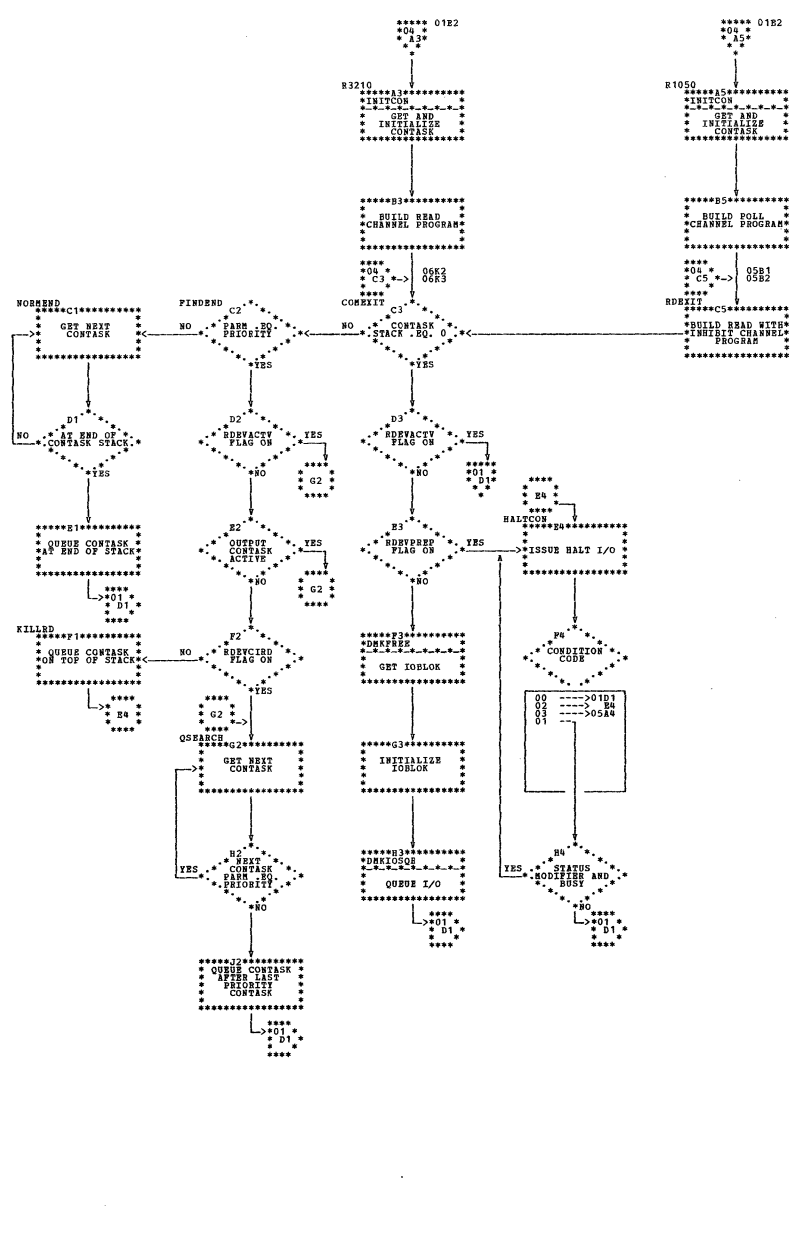
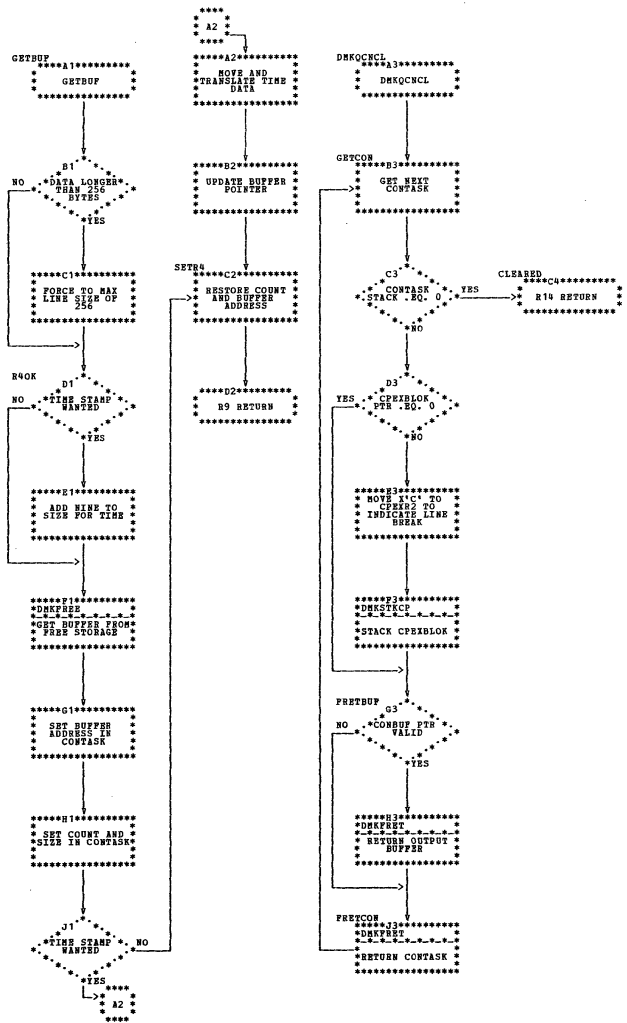
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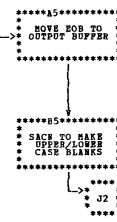
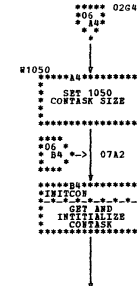
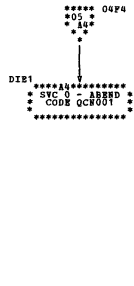
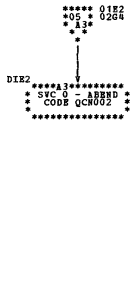
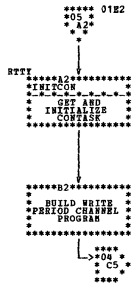
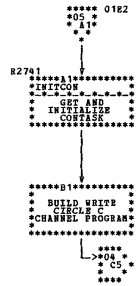
DMKPTR -- Real Storage Page Manager (Part 13 of 13)



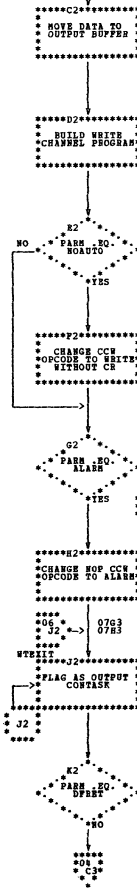


DMKQCN -- Console Message Queue Manager (Parts 3 and 4 of 7)

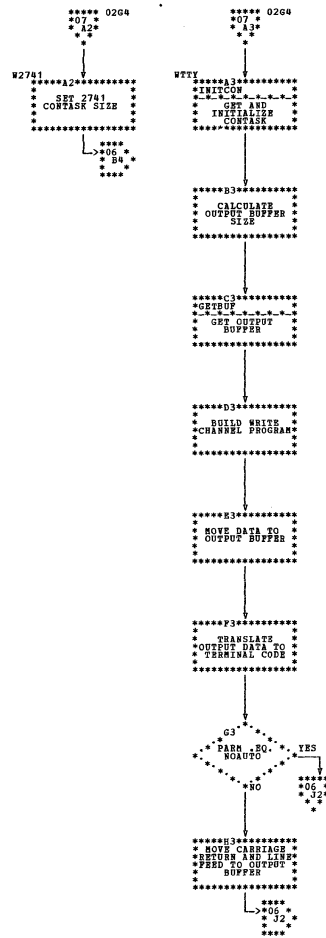




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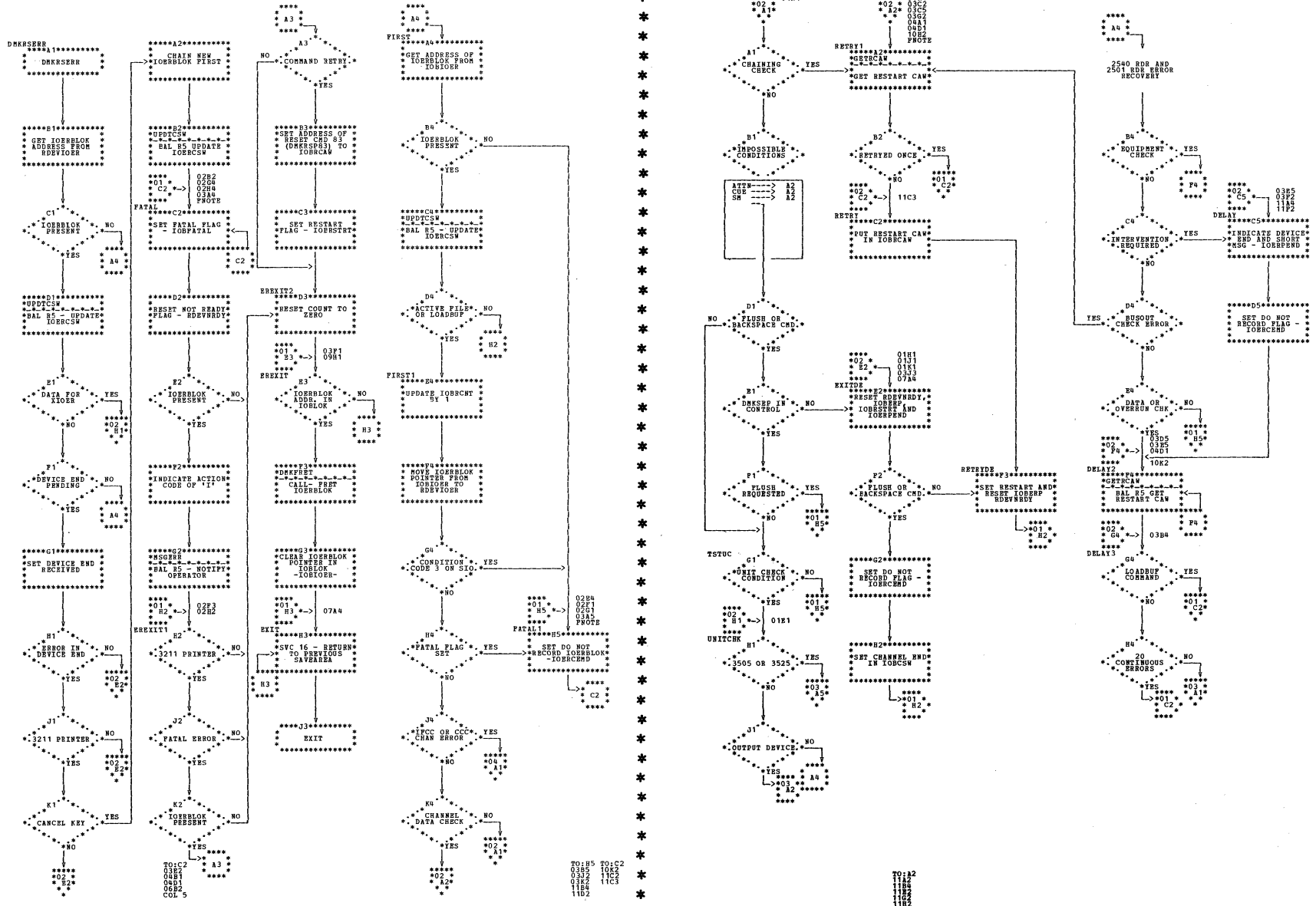


DMKQCN -- Console Message Queue Manager (Part 7 of 7)

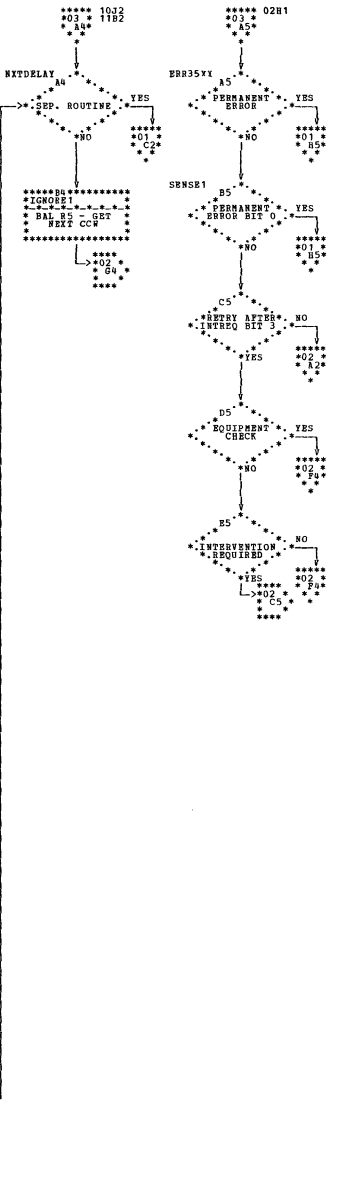
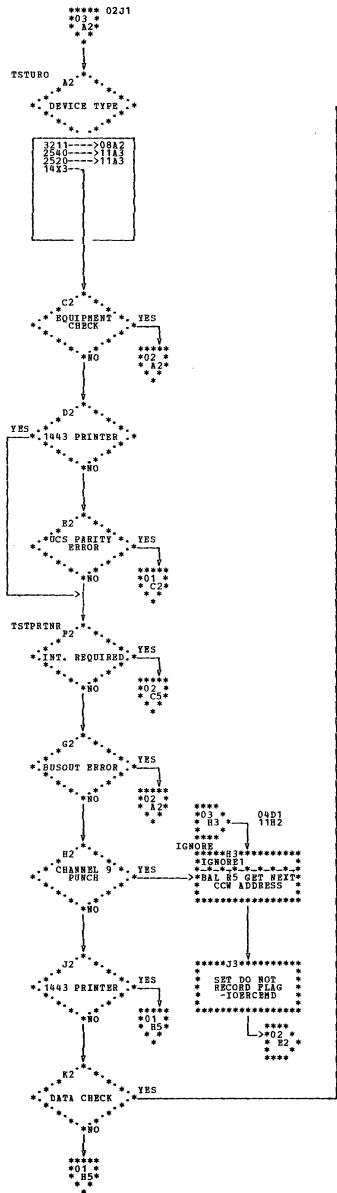
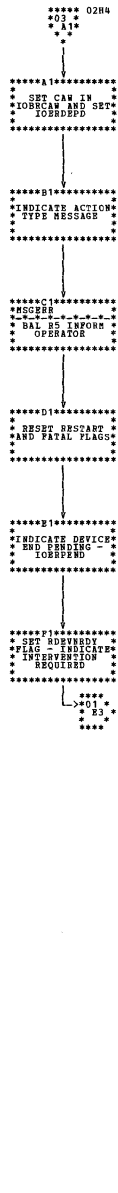




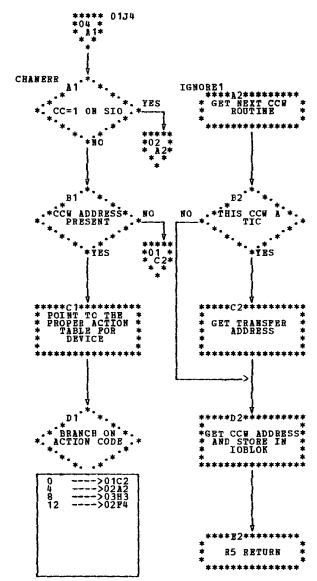
DMKRSE -- Real Spool Error Procedures (Parts 1 and 2 of 11)



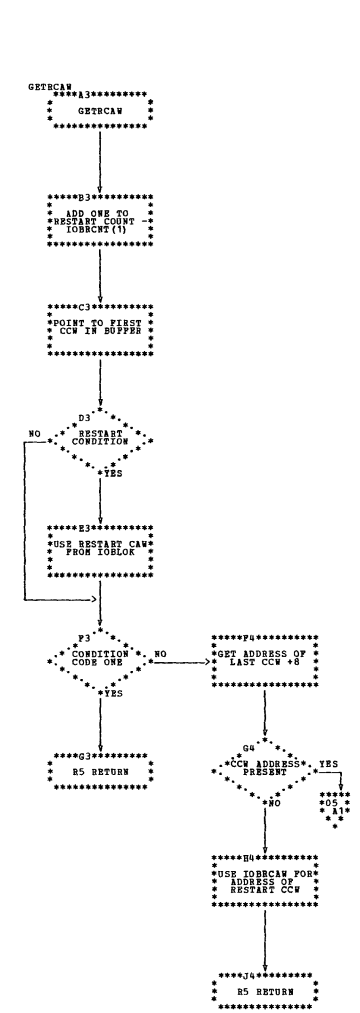




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0	>01C2
8	>02Z4
9	>03B3
12	>02F4

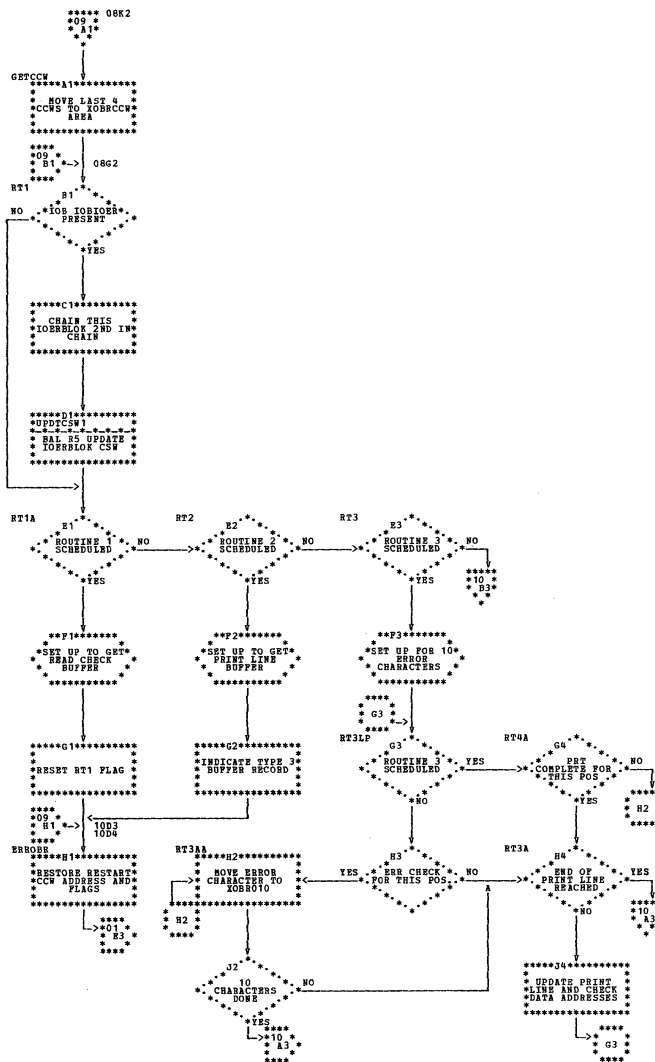


| DMKRSE -- Real Spool Error Procedures (Parts 3 and 4 of 11)

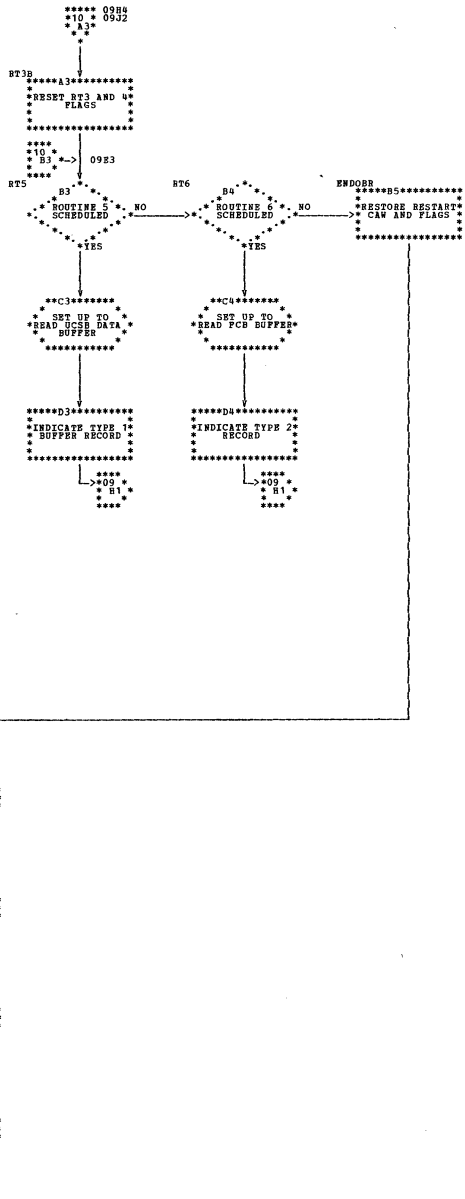
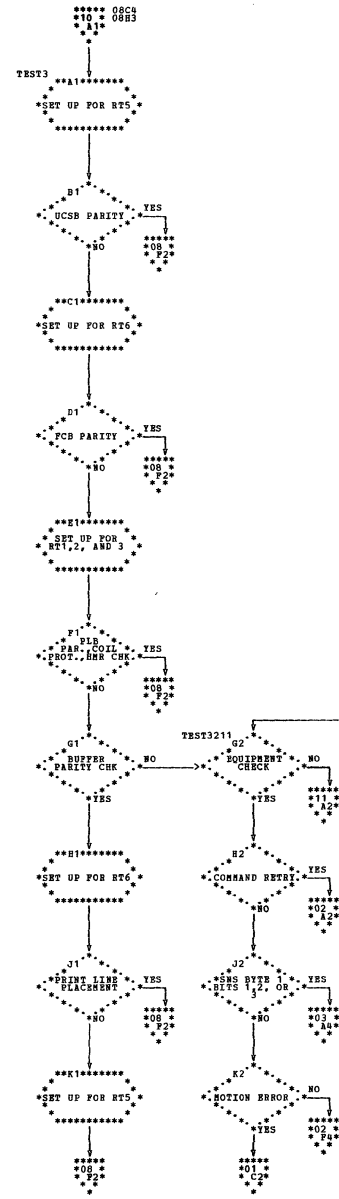


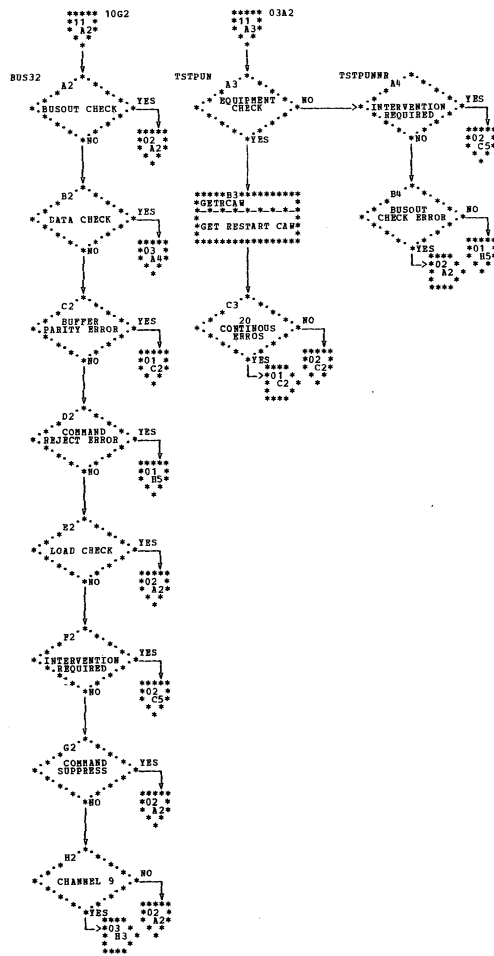


| DMKRSE -- Real Spool Error Procedures (Parts 9 and 10 of 11)



\* \* \* \* \*





| DMKRSE -- Real Spool Error Procedures (Part 11 of 11)



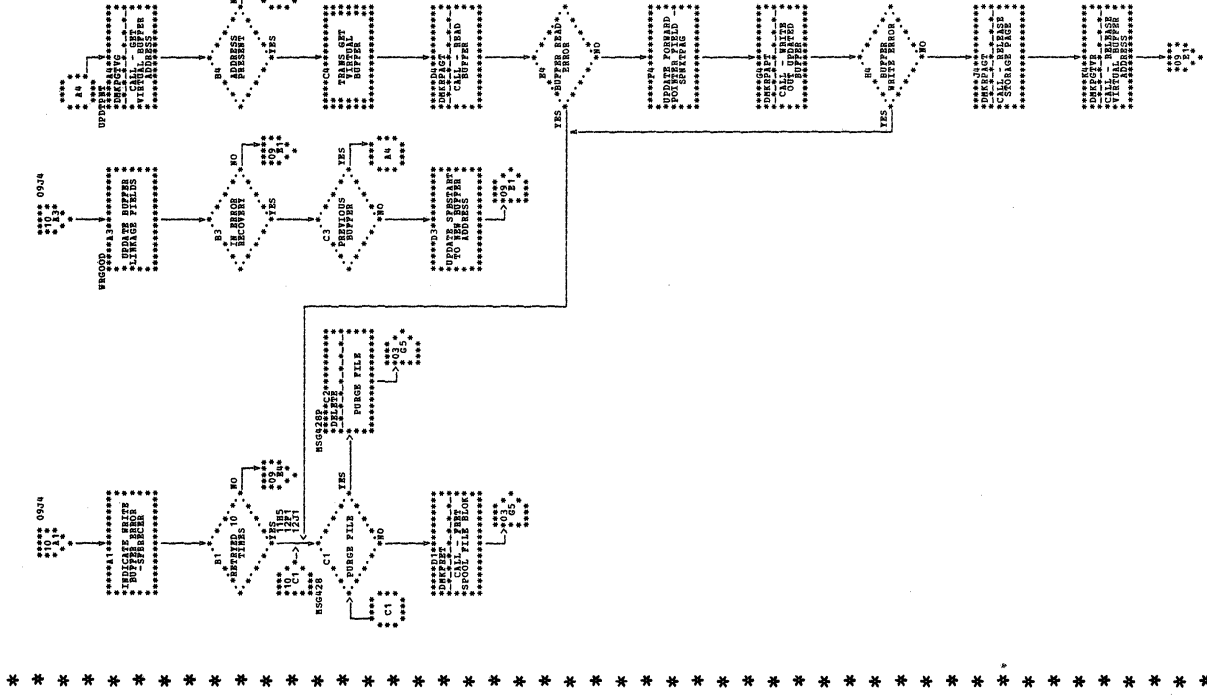




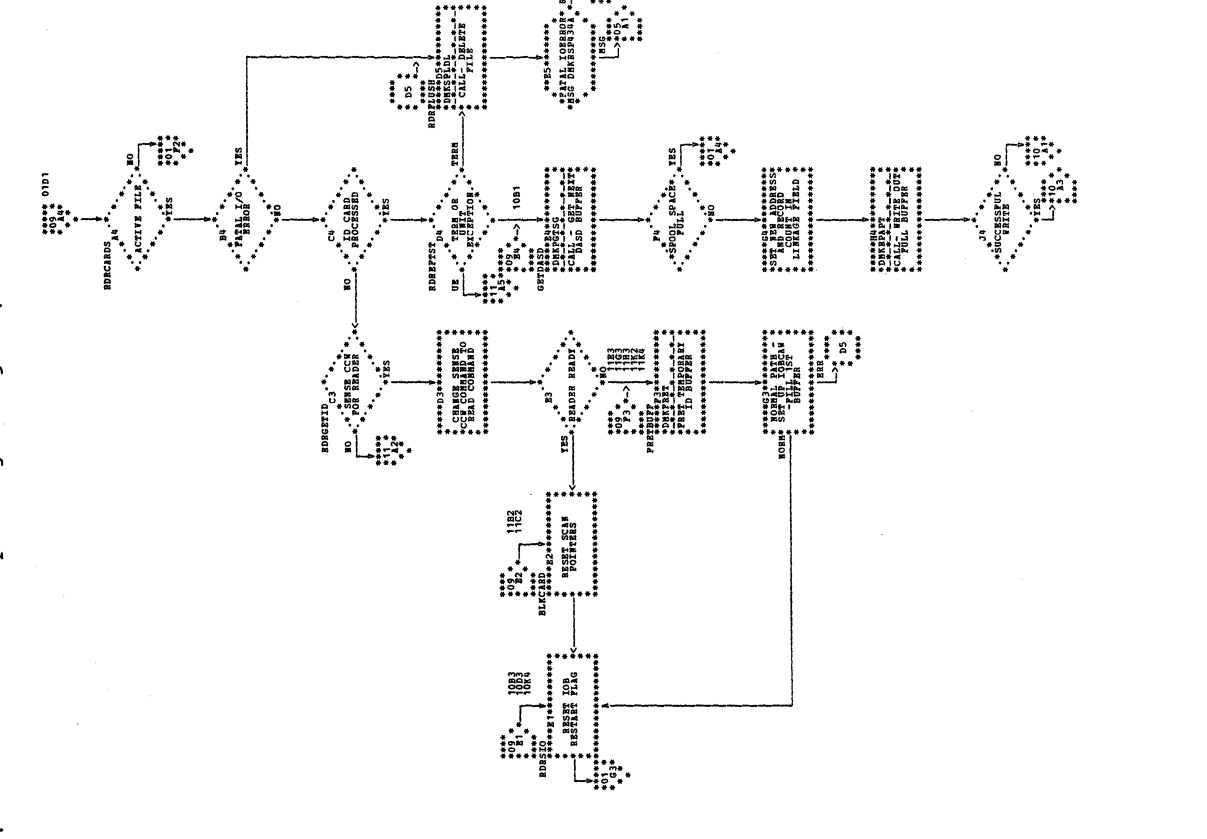


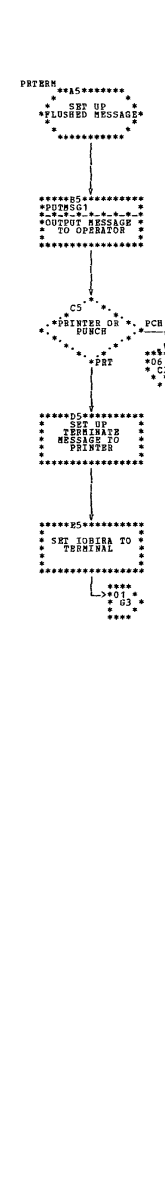
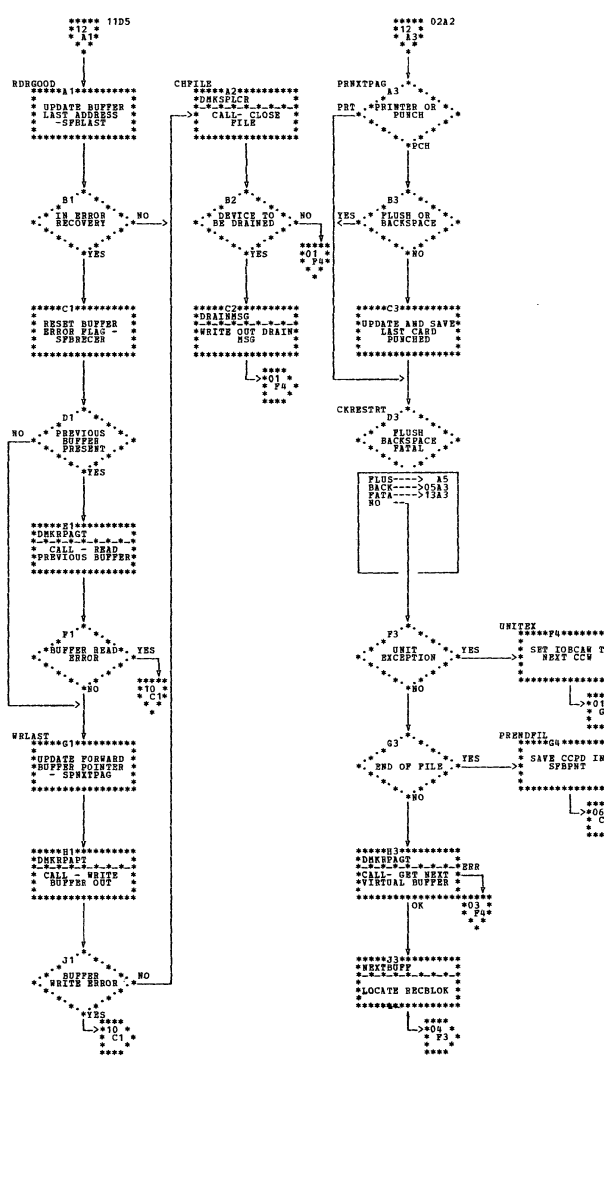
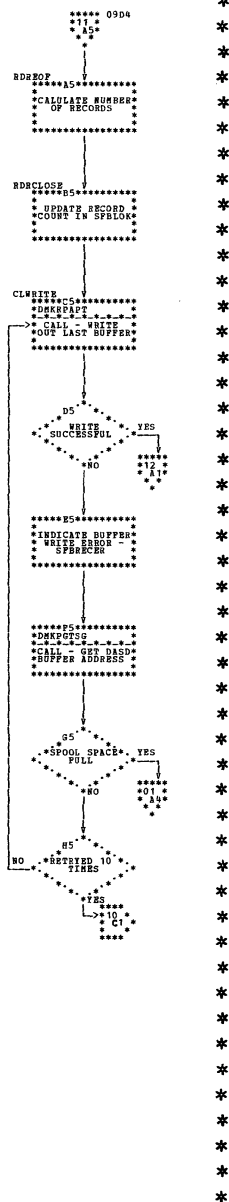
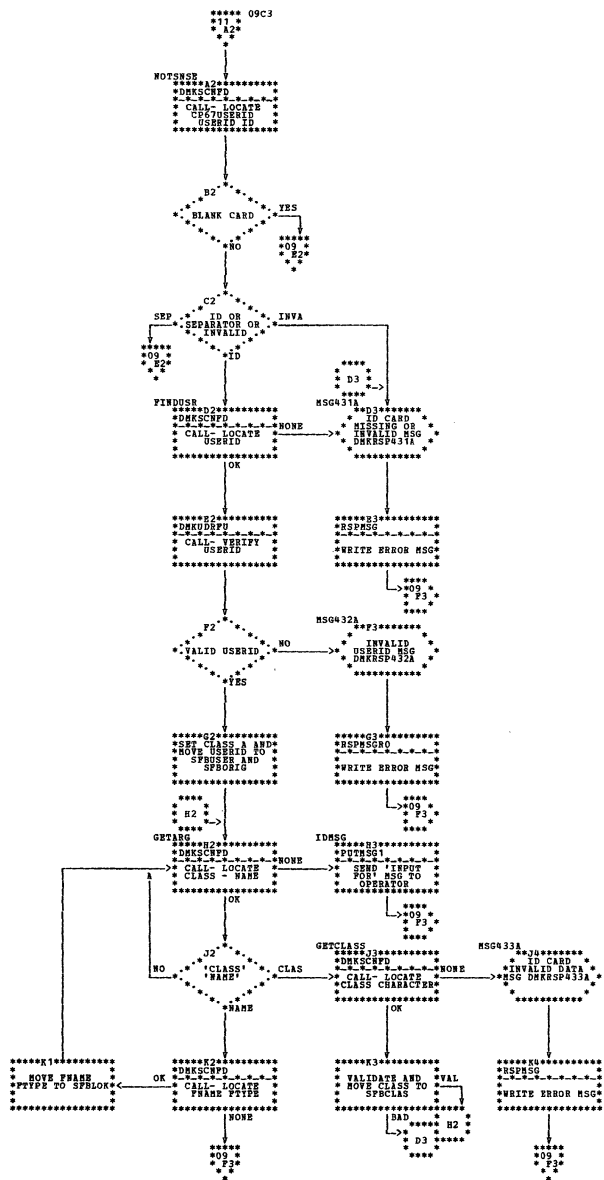


DMKRSP -- Real Spooling Manager (Parts 9 and 10 of 13)

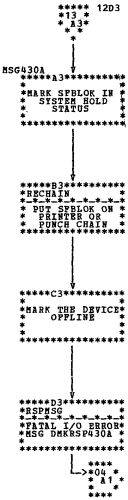


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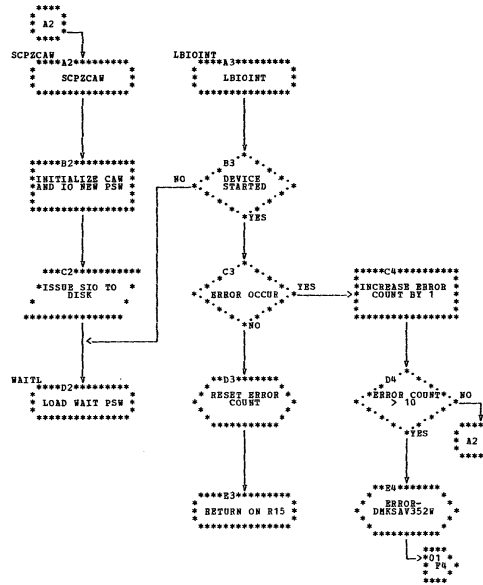


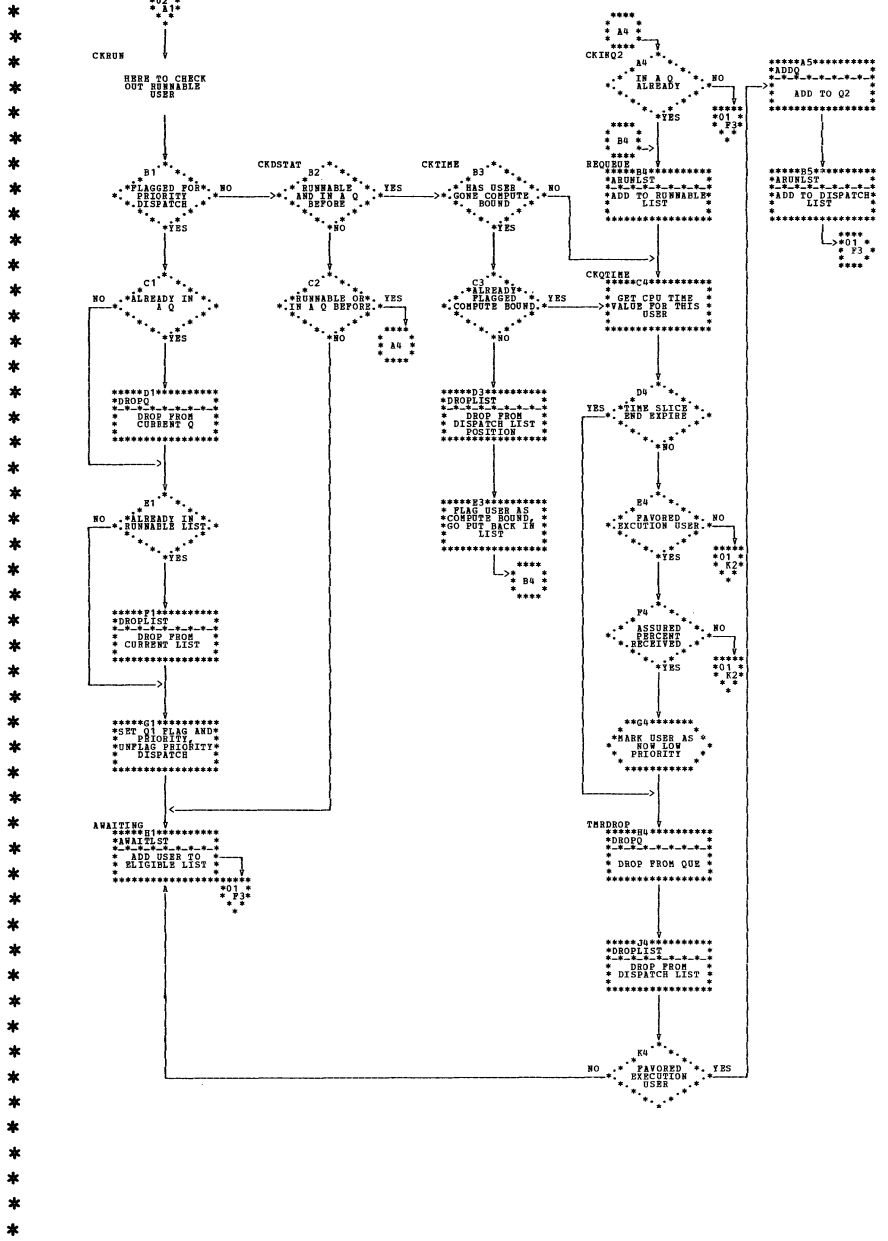
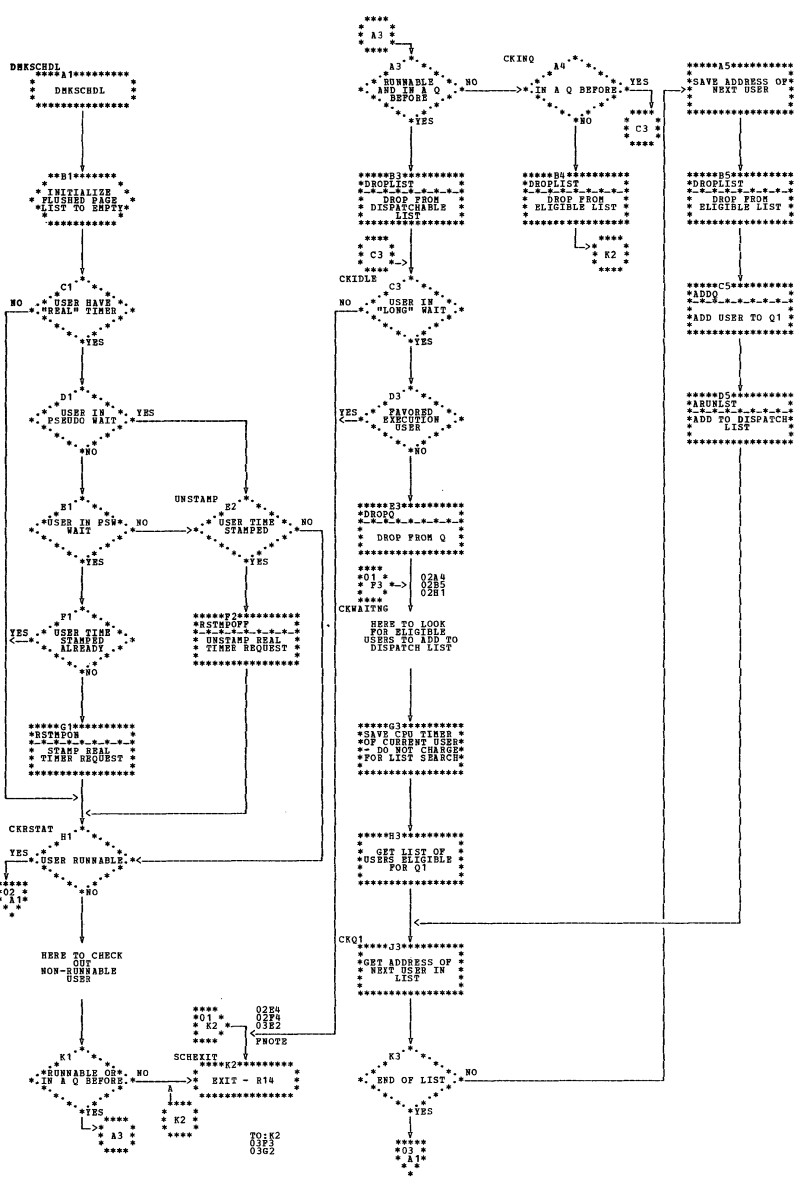
| DMKRSP -- Real Spooling Manager (Part 13 of 13)





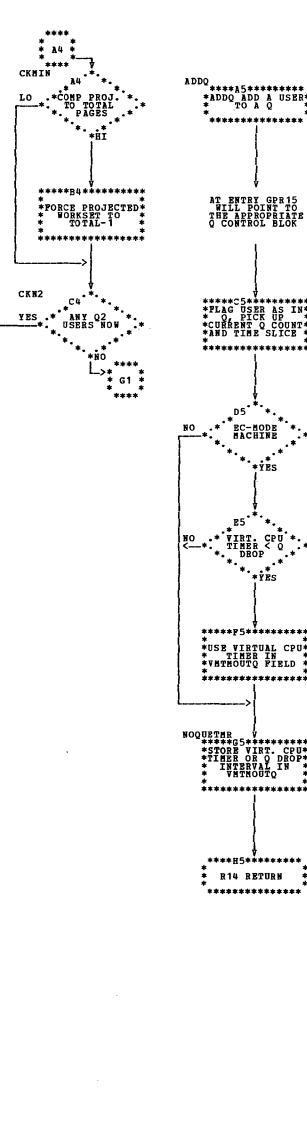
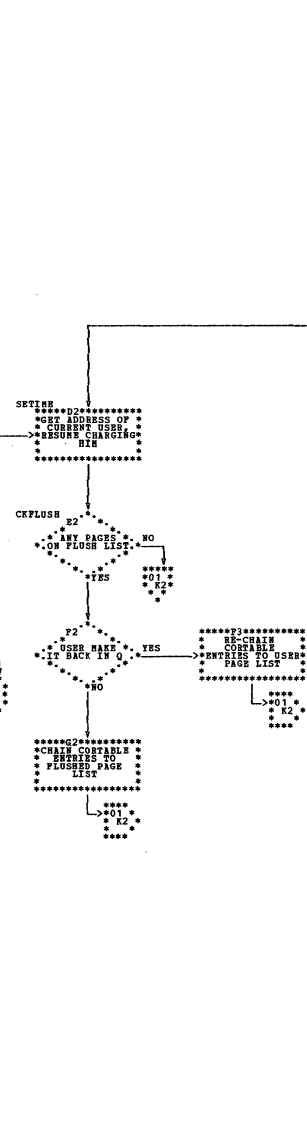
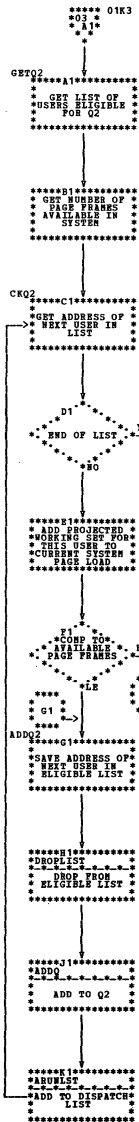
DMKSAV -- Save CP Nucleus or SYSRES (Part 3 of 3)



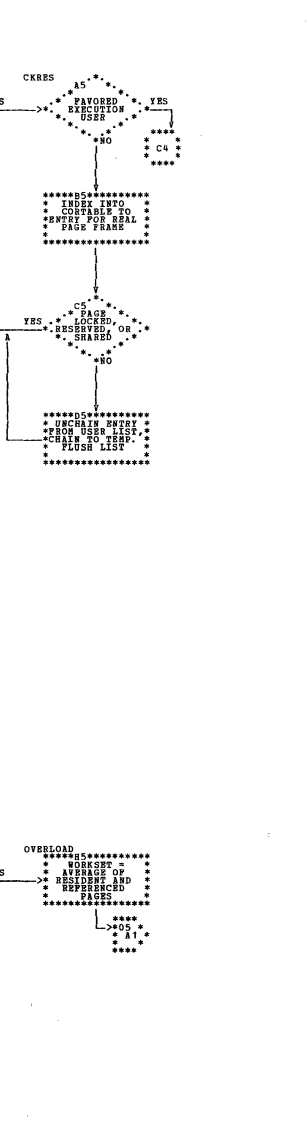
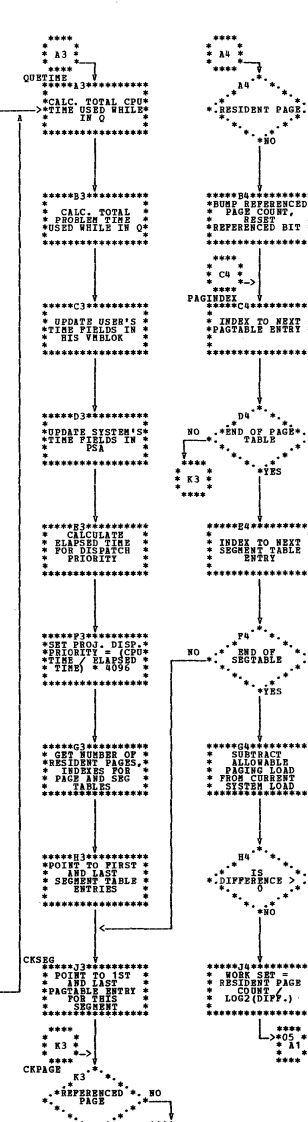
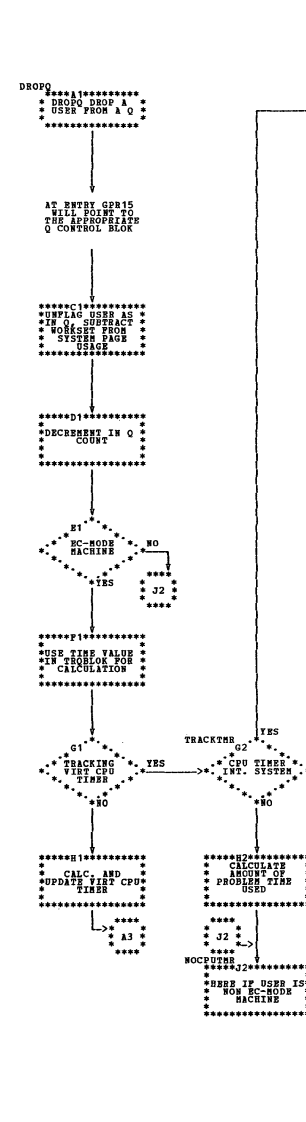


| DMKSCH -- Scheduler (Parts 1 and 2 of 8)

DMKSCH -- Scheduler (Parts 3 and 4 of 8)



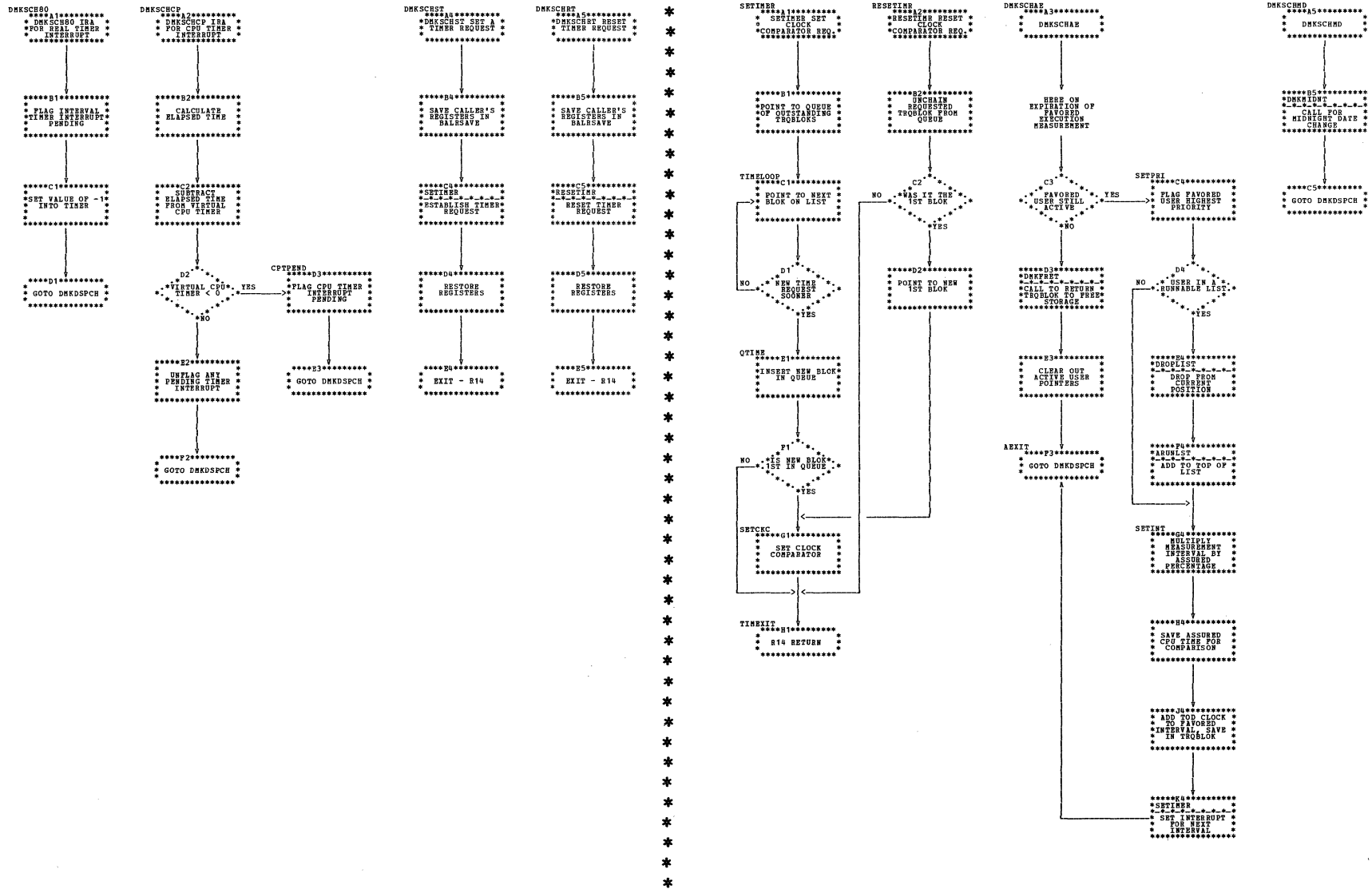
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| DMKSCH -- Scheduler (Parts 7 and 8 of 8)



DMKSCNRU  
\*\*\*\*\*1\*\*\*\*\*  
DMKSCNRU

FIND THE REAL CHANNEL CONTROL UNIT, AND DEVICE BLOCKS

IF SUCCESSFUL ON FINDING GPR-9=CHANNEL GPR-7=CONTROL UNIT GPR-8=DEVICE

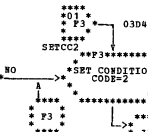
\*\*\*\*\*D1\*\*\*\*\*  
INDEX INTO CHANNEL TABLE

\*CHANNEL EXIST\*  
NO

\*\*\*\*\*P1\*\*\*\*\*  
RETURN TO CALLER

\*\*\*\*\*E2\*\*\*\*\*  
INDEX INTO CONTROL UNIT TABLE

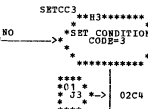
\*CONTROL UNIT EXISTS\*  
NO



\*\*\*\*\*G2\*\*\*\*\*  
INDEX INTO DEVICE TABLE

\*DEVICE EXIST\*  
NO

\*\*\*\*\*S2\*\*\*\*\*  
RETURN TO CALLER



DMKSCVU  
\*\*\*\*\*1\*\*\*\*\*  
DMKSCVU

FIND THE BLOCKS FOR A GIVEN VIRTUAL DEVICE ADDRESS

IF SUCCESSFUL ON FINDING GPR-9=VCHNLOK GPR-7=VDEVLOK GPR-8=VDEVELOK

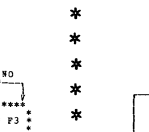
\*\*\*\*\*D4\*\*\*\*\*  
INDEX INTO VIRTUAL CHANNEL TABLE

\*CHANNEL EXIST\*  
NO

\*\*\*\*\*P4\*\*\*\*\*  
RETURN TO CALLER

\*\*\*\*\*E5\*\*\*\*\*  
INDEX INTO VIRTUAL CONTROL UNIT TABLE

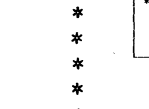
\*CONTROL UNIT EXISTS\*  
NO



\*\*\*\*\*G5\*\*\*\*\*  
INDEX INTO VIRTUAL DEVICE TABLE

\*VIRTUAL DEVICE EXIST\*  
NO

\*\*\*\*\*S5\*\*\*\*\*  
RETURN TO CALLER



\* \* \* \* \*

DMKSCVWS  
\*\*\*\*\*1\*\*\*\*\*  
DMKSCVWS

SEARCH FOR A DEVICE WITH A GIVEN SERIAL NUMBER

EXIT CONDITION CCO=SEARCH, CCI=NO SEARCH, CCB=CELL=ERROR

\*\*\*\*\*D1\*\*\*\*\*  
LOAD ADDRESS OF FIRST DEVVELOK

\*GET NUMBER OF DEVVELOK IN THE SYSTEM\*

VSERLOOP P1  
\*APPLIES OR DISAPPLIES TO SYSTEM\*

\*ATTACHED TO SYSTEM\*

\*ALL DEVVELOKS SEARCHED\*

\*\*\*\*\*J1\*\*\*\*\*  
SET CONDITION CODE=1

VSERWEXIT  
RETURN TO CALLER

DMKSCVWD  
\*\*\*\*\*1\*\*\*\*\*  
DMKSCVWD

COMPUTE REAL DEVICE ADDRESS IN CCU FORM

ON EXISTING GPR-1=CCU

\*\*\*\*\*D2\*\*\*\*\*  
OF CHANNEL, GET LUID, & DEVICE ADDR, INTO GPR-1

\*\*\*\*\*P2\*\*\*\*\*  
RETURN TO CALLER

DMKSCVU  
\*\*\*\*\*3\*\*\*\*\*  
DMKSCVU

CONVERT VIRTUAL DEVICE ADDRESS IN CCU FORM

ON EXISTING GPR-1=VIRTUAL CCU

\*\*\*\*\*D3\*\*\*\*\*  
GET VIRTUAL CHANNEL ADDRESS

\*CHANNEL EXISTS\*  
NO

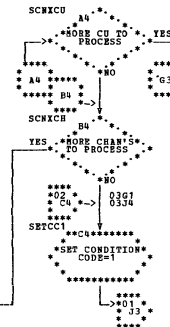
\*\*\*\*\*P3\*\*\*\*\*  
POINT TO VIRTUAL CONTROL UNIT

\*\*\*\*\*G3\*\*\*\*\*  
CONTROL UNIT EXIST

\*\*\*\*\*H3\*\*\*\*\*  
POINT TO VIRTUAL DEVICE

\*DEVICE EXIST\*  
NO

\*THE ONE WE'RE LOOKING FOR\*



\*\*\*\*\*5\*\*\*\*\*  
OF CHANNEL GET LUID, & DEVICE ADDR, INTO GPR-1.  
RETURN TO CALLER

DMKSCN -- Scan Routines (Parts 3 and 4 of 4)

DMKSCN A1

FIND THE VNBLOK WITH A GIVEN SSERID  
IF SUCCESSFUL OR EXISTING GPR-1 POINTS TO VNBLOK

D1 START WITH SYSTEMS VNBLOK

FINDLOOP  
USERID MATCH

F1 GET NEXT VNBLOK

G1 ALL VNBLOKS SEARCHED

DMKSCNPD A3

LOCATE NEXT FIELD IN THE INPUT BUFFER  
EXIT-CONDITION C1=FIELD END C2=CARET C3=END BUFFER

D3 BUFFER COUNT ZERO

SCANPRST ES BLANK CHARACTER

F3 POINT TO NEXT BUFFER POSITION

SCANZERD D4 RESTORE REGISTERS

SCANSTRT E2 SAVE POINTER TO START OF FIELD

SCANLAST F4 END OF FIELD

G4 LOGICAL LIVE END CHAR

SCANLEND H5 SAVE ADDRESS OF NEXT BUFFER FIELD

J4 INPUT FIELD FOUND

DMKSCNLI A3

FIND ALL THE LINKS TO A GIVEN MINIDISK  
EXIT-CONDITION C1=NO LINKS C2=NO DISK C3=WRITE DISKS

D5 GET ADDRESS OF FIRST VIRTUAL DEVICE BLOCK

LINKCLR E5 ANY VDEVLOCKS FOR THIS USER

LINKKCD F5 VALID VIRTUAL VNBLOK

G5 VNBLOK NOVNBLOK

LINKRVT H5 GET NEXT VIRTUAL DEVICE BLOCK

J5 ANY MORE LEFT

\* \* \* \* \*

LINKWDR A1 GET NEXT USERS VNBLOK

B1 END OF VNBLOK CHAIN

LINKREREC I SET NUMBER OF LINKS IN GPR-1 SET COMPILATOR CODE

D1 RETURN TO CALLER

LINKKIND A2 DISK EXTENT MATCH

B2 DEDICATED DEVICE

C2 TYPE PSEUDO A 2311

D2 REAL DEVICE A 2311

LINKKIND E2 SAVE ADDRESS OF FIRST LINKED VDEVLOCK

LINKNEXT F2 R-O DISK

LINKCOM H2 ALL VDEVLOCKS SEARCHED

LINKREREC I SET NUMBER OF LINKS IN GPR-1 SET COMPILATOR CODE

D2 REAL DEVICE A 2311

LINKKIND E2 SAVE ADDRESS OF FIRST LINKED VDEVLOCK

LINKNEXT F2 R-O DISK

LINKCOM H2 ALL VDEVLOCKS SEARCHED

DMKSCNRE A4 FIND THE DEVICE NAME FOR A GIVEN DEVICE ADDRESS

EXIT-CONDITION GPR-1 NAME IN BRIDGE OF DEVICE

D4 GET DEVICE CLASS FROM VNBLOK - TOVNBLOK

E4 DEVELOP INDEX VALUE FROM DEVICE CLASS

F4 LOAD GPR-1 WITH DEVICE NAME

G4 RETURN TO CALLER

D4 GET DEVICE CLASS FROM VNBLOK - TOVNBLOK

E4 DEVELOP INDEX VALUE FROM DEVICE CLASS

F4 LOAD GPR-1 WITH DEVICE NAME

G4 RETURN TO CALLER

DMKSCNRE A4 FIND THE DEVICE NAME FOR A GIVEN DEVICE ADDRESS

EXIT-CONDITION GPR-1 NAME IN BRIDGE OF DEVICE

D4 GET DEVICE CLASS FROM VNBLOK - TOVNBLOK

E4 DEVELOP INDEX VALUE FROM DEVICE CLASS

F4 LOAD GPR-1 WITH DEVICE NAME

G4 RETURN TO CALLER

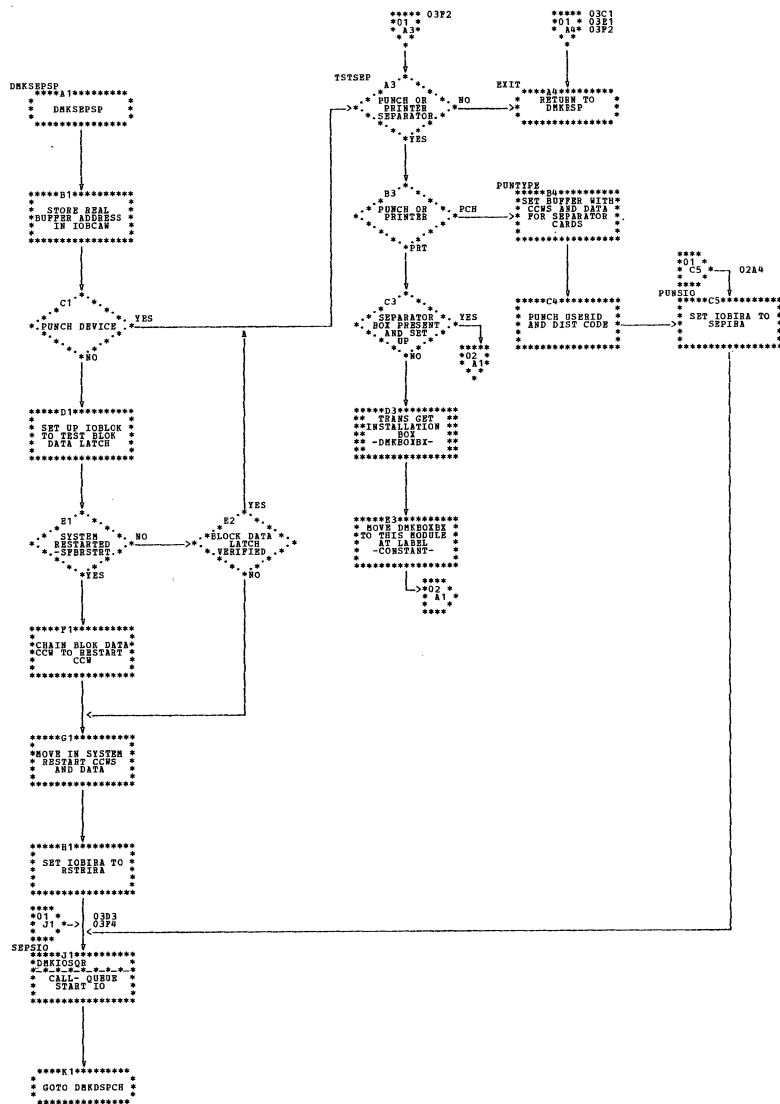
D4 GET DEVICE CLASS FROM VNBLOK - TOVNBLOK

E4 DEVELOP INDEX VALUE FROM DEVICE CLASS

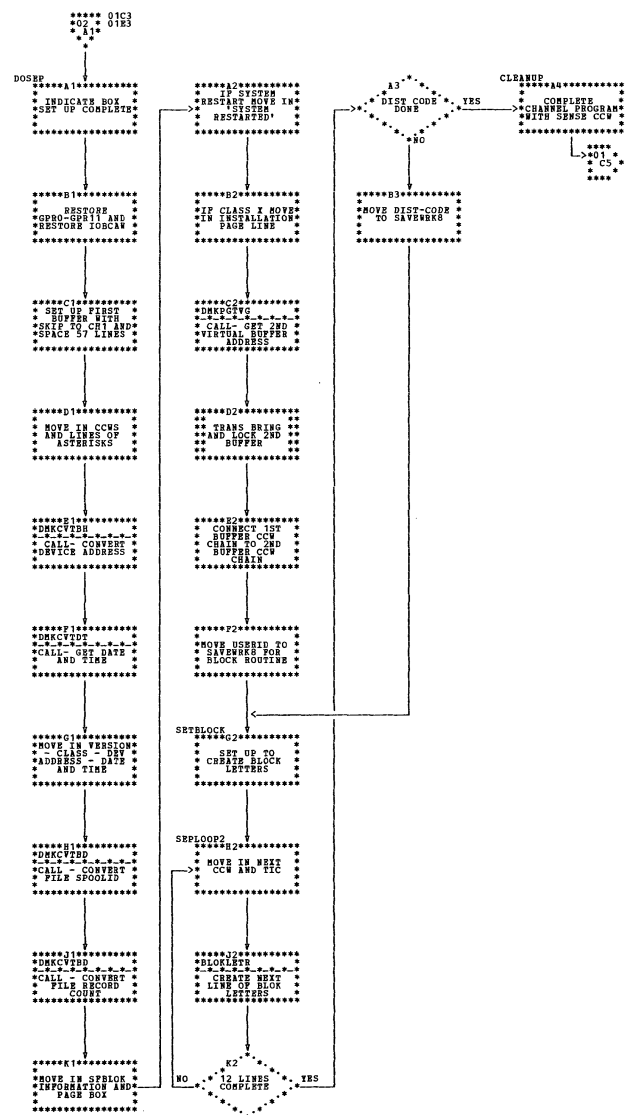
F4 LOAD GPR-1 WITH DEVICE NAME

G4 RETURN TO CALLER

\* \* \* \* \*

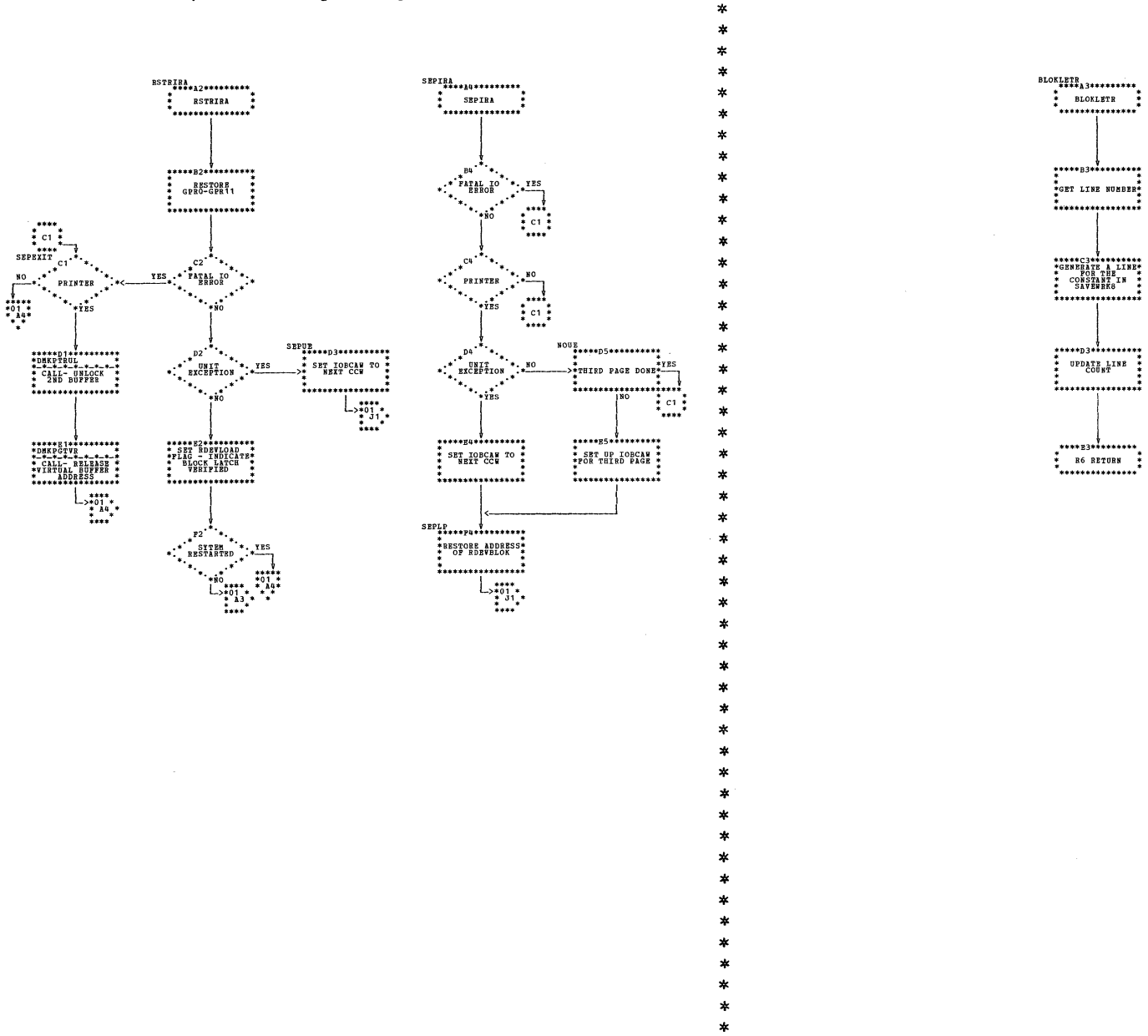


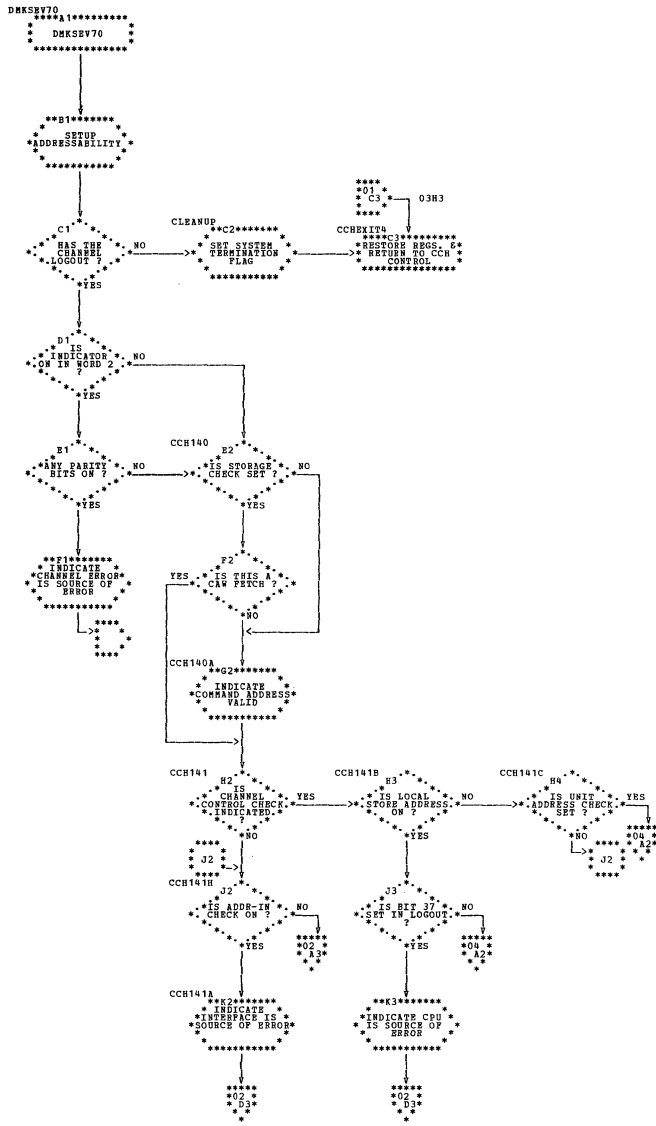
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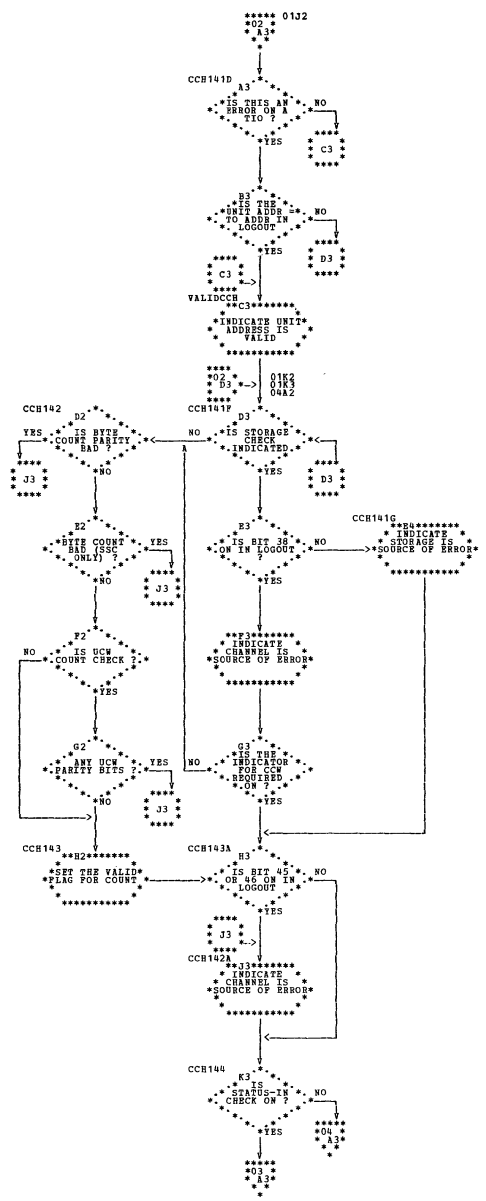
DMKSEP -- Print/Punch Output Separator (Parts 1 and 2 of 4)

DMKSEP -- Print/Punch Output Separator (Parts 3 and 4 of 4)





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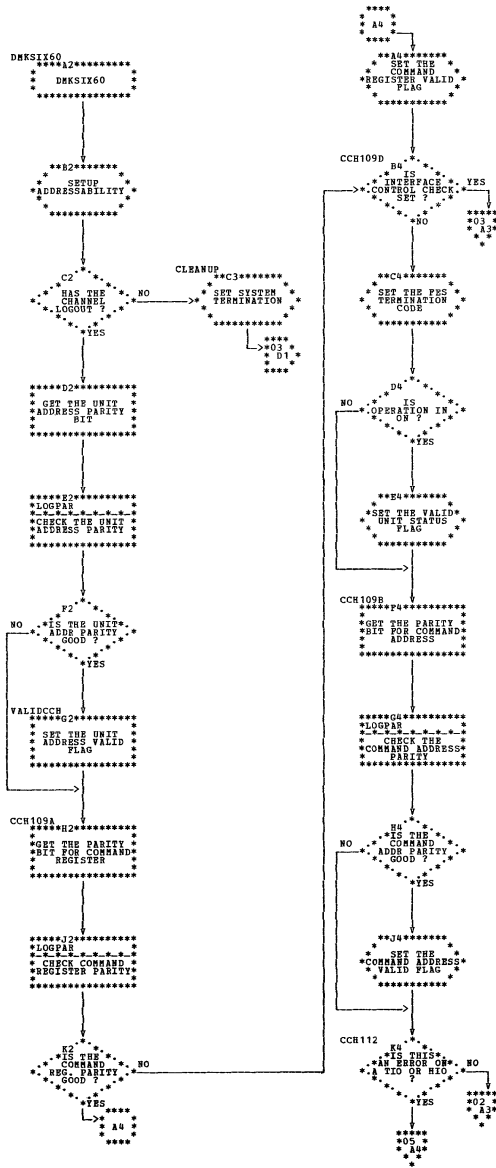
| DMKSEV -- 2870 Channel Module (Parts 1 and 2 of 6)



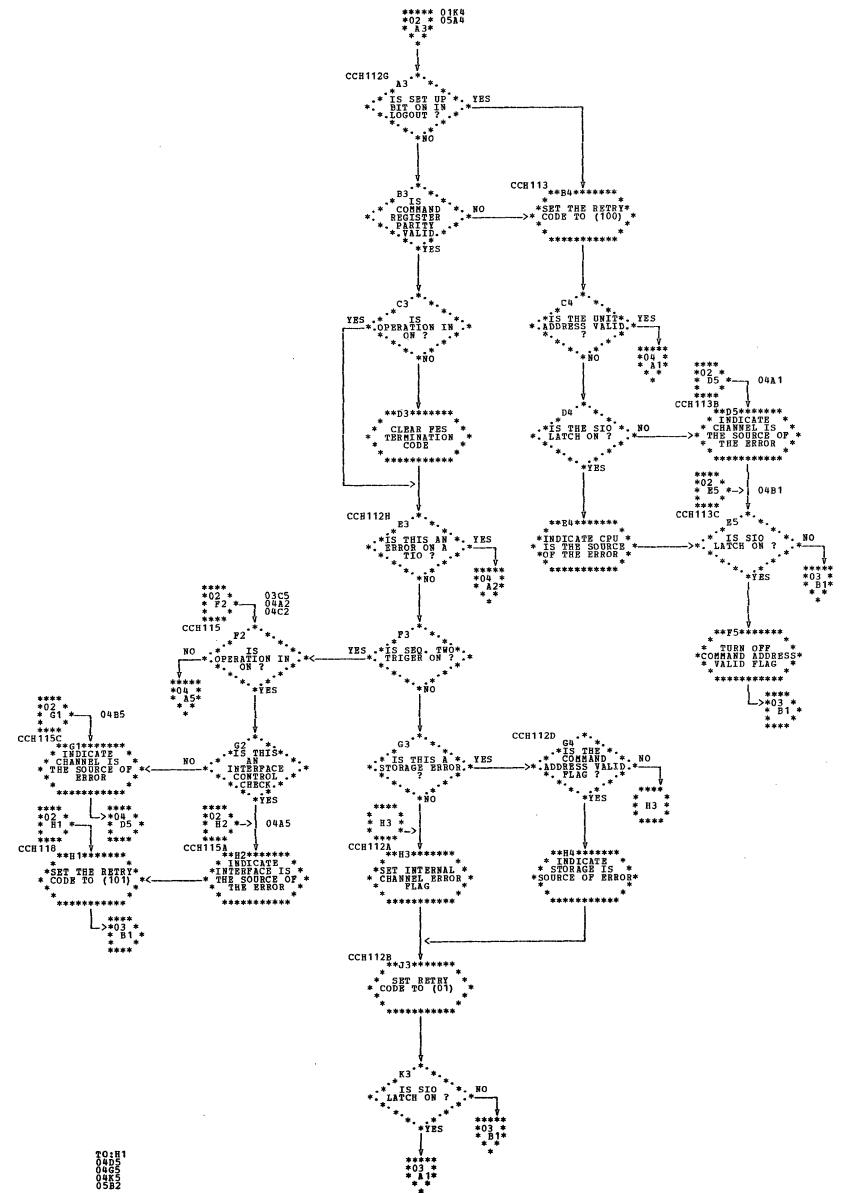




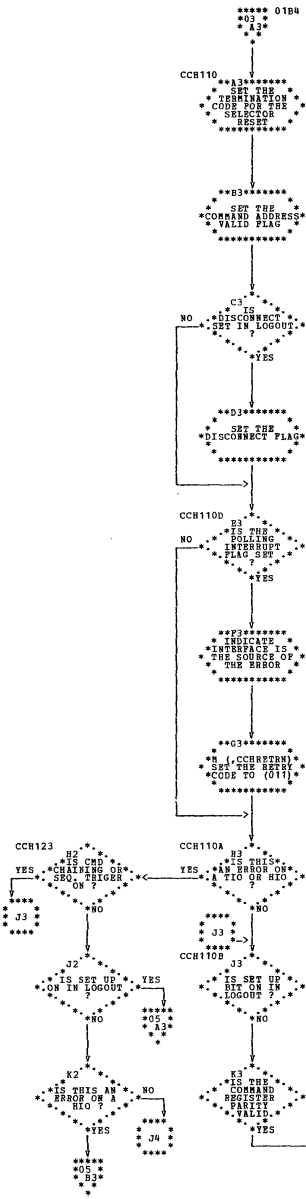
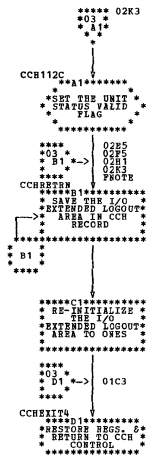
DMKSIX -- 2860 Channel Module (Parts 1 and 2 of 5)



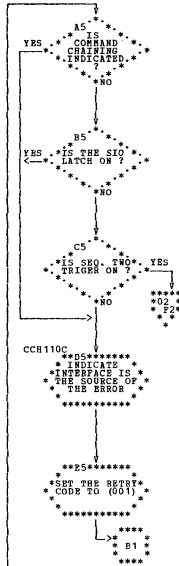
\* \* \* \* \*



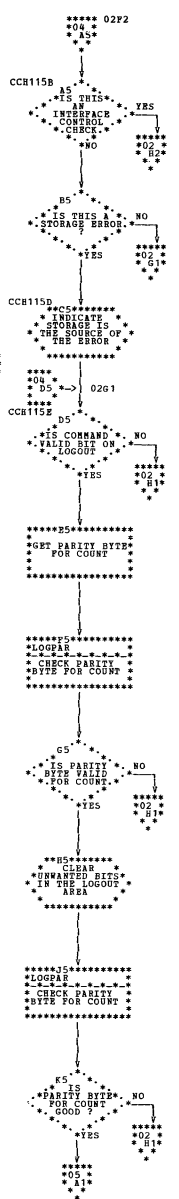
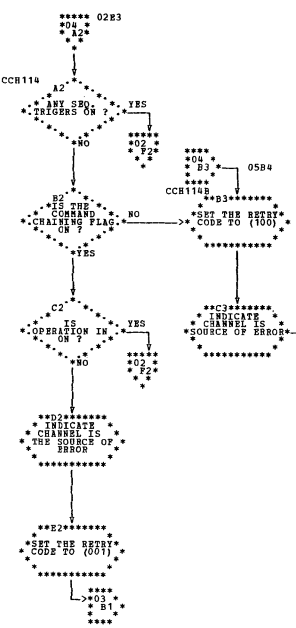
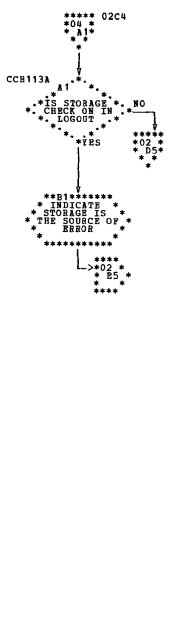
O1B1  
O1C1  
O1D1  
O1E1



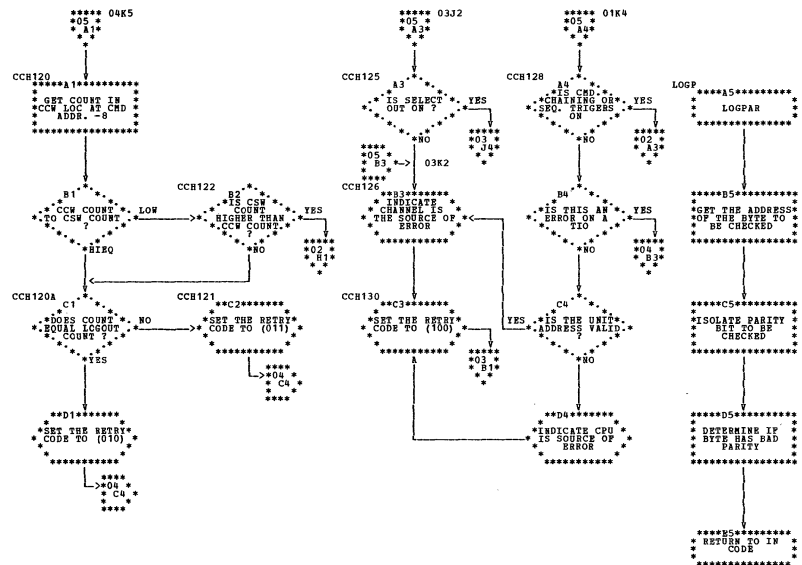
FO:B1  
02K4  
02K5  
02K6  
02K7  
02K8



\* \* \* \* \*

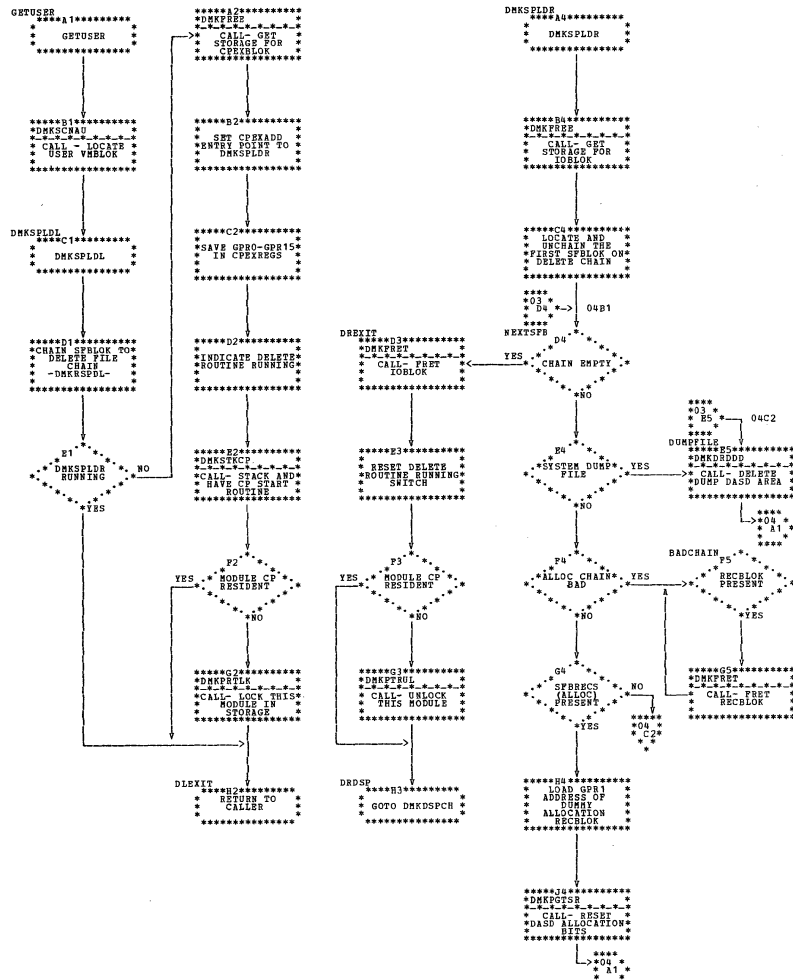


( DMKSIX -- 2860 Channel Module (Part 5 of 5)

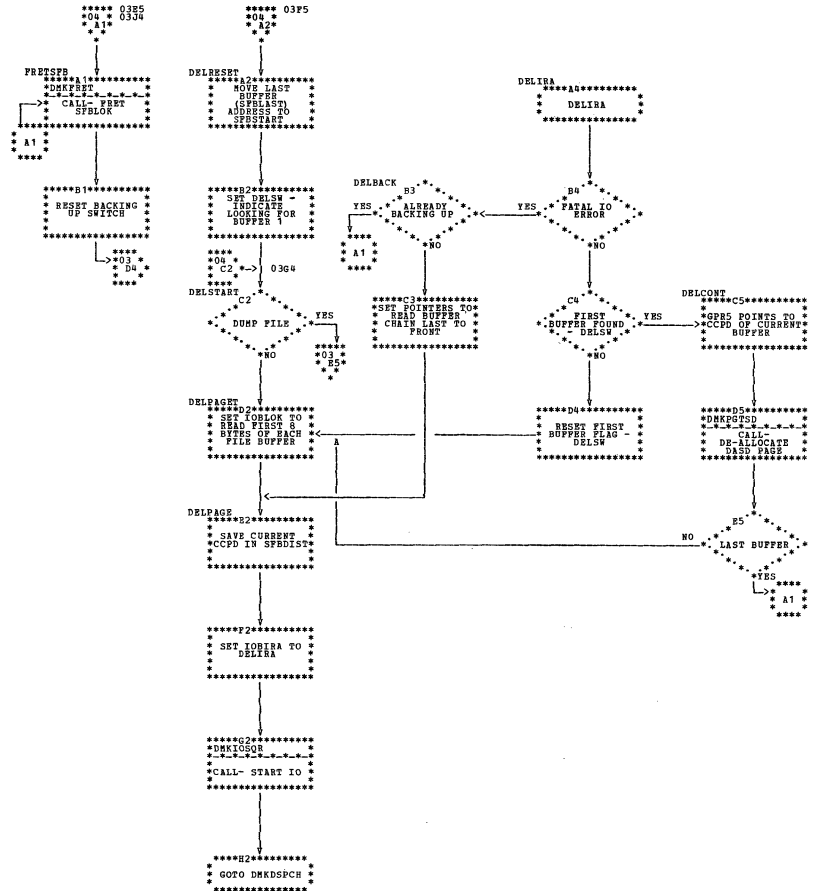




DMKSPL -- Spooling Subroutines (Parts 3 and 4 of 5)



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```
GETID .....
*****
ASSIGN SPOOLID
*****
|
*****
UPDATE SPOOLID
  BB IF 9900
  RESTART AT 1
*****
|
*****
RETURN TO
  CALLER
*****
```

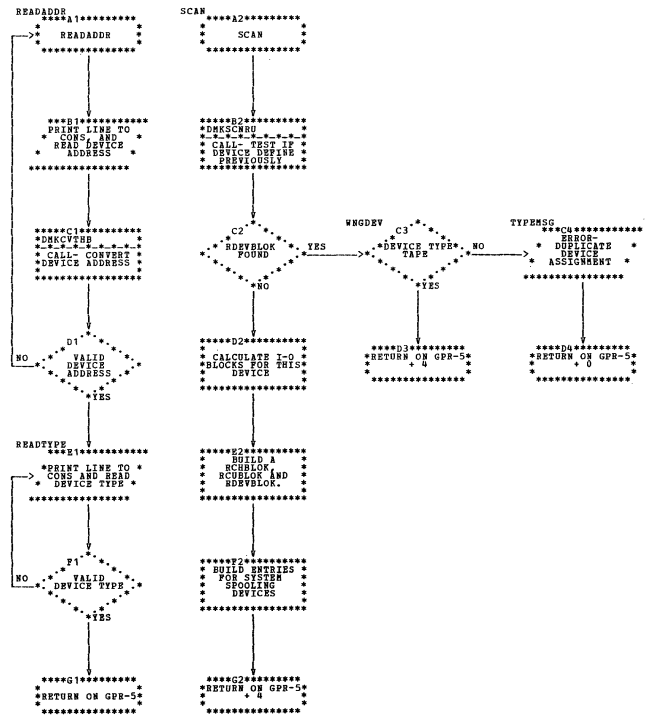
| DMKSPL -- Spooling Subroutines (Part 5 of 5)

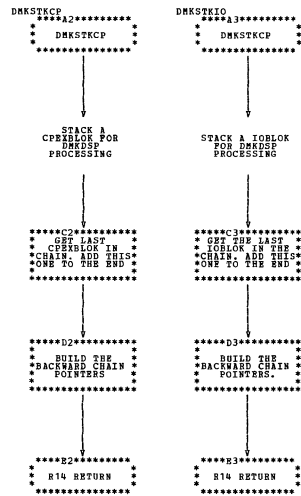






DMKSSP -- Build Real I/O Blocks for Starter System or at IPL Time (Part 3 of 3)



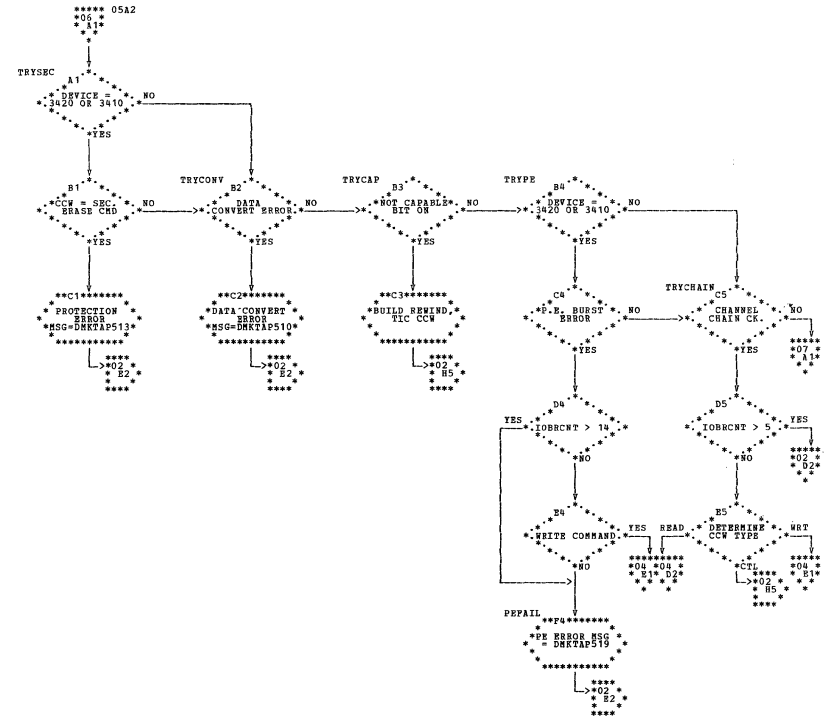
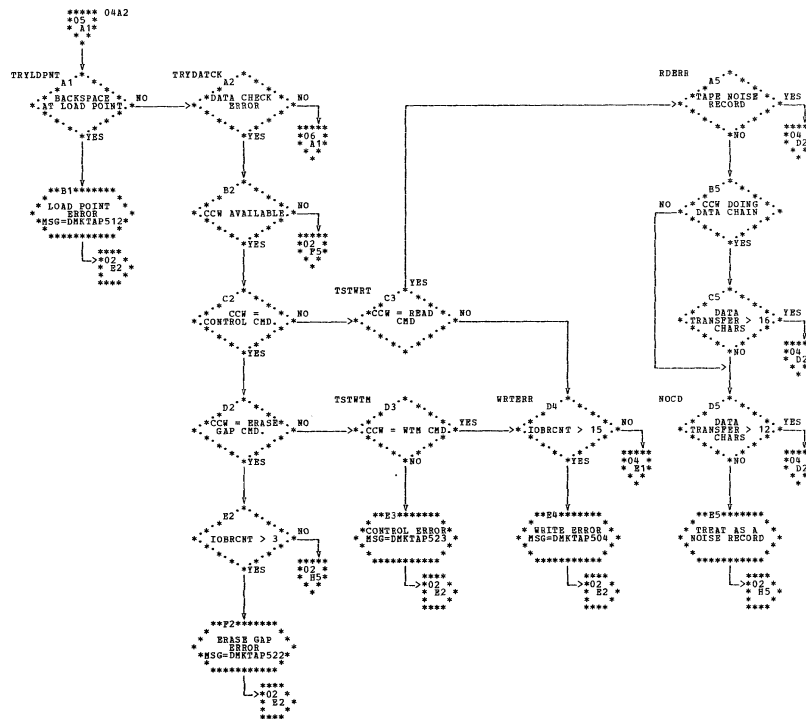


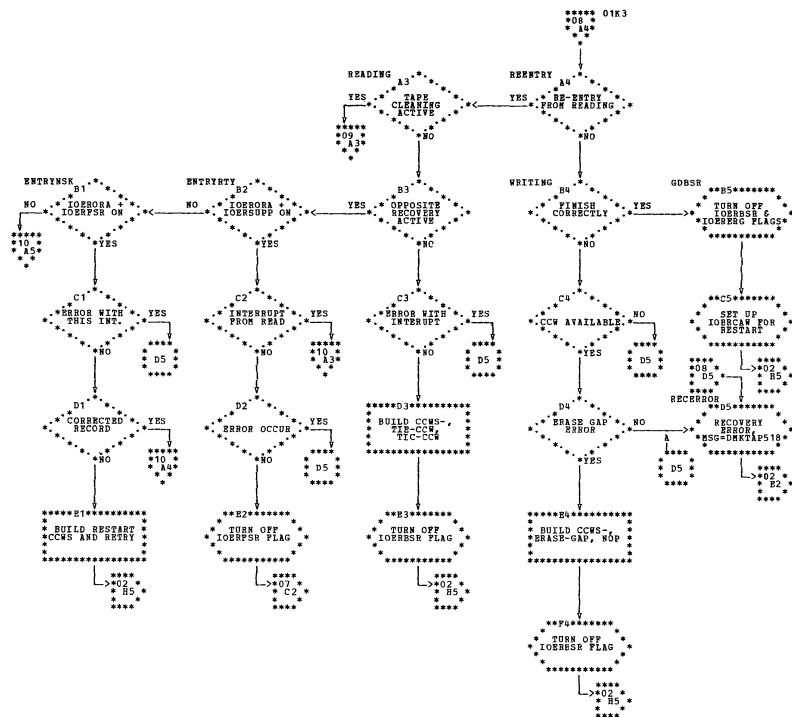
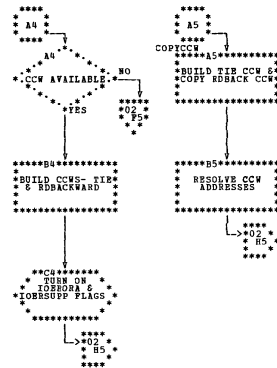
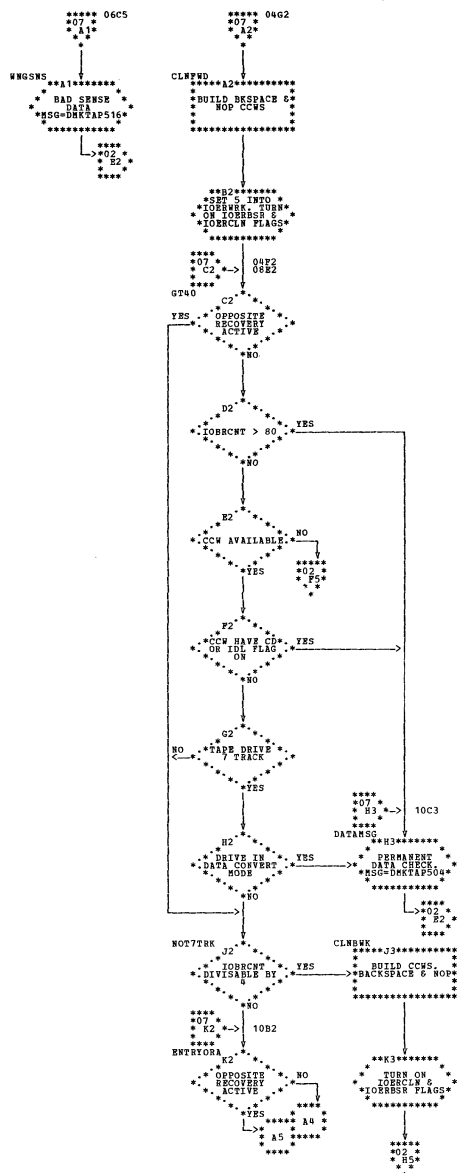
DMKSTK -- Stack (queue) a CPEXBLOK or IOBLOK for Dispatching (Part 1 of 1)





DMKTAP -- Tape Error Recovery Procedures (Parts 5 and 6 of 11)



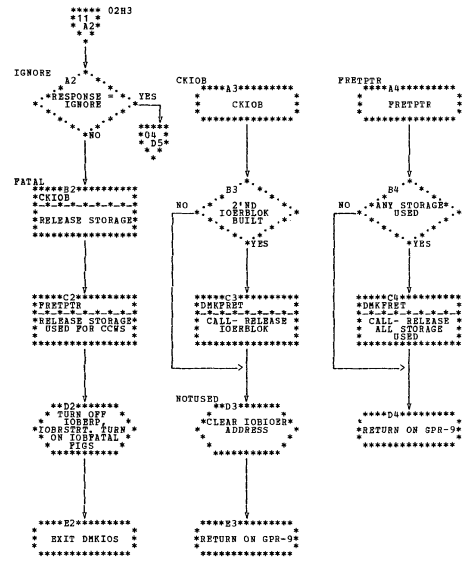


| DMKTAP -- Tape Error Recovery Procedures (Part 7 and 8 of 10)

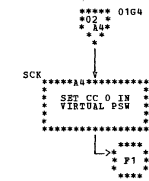
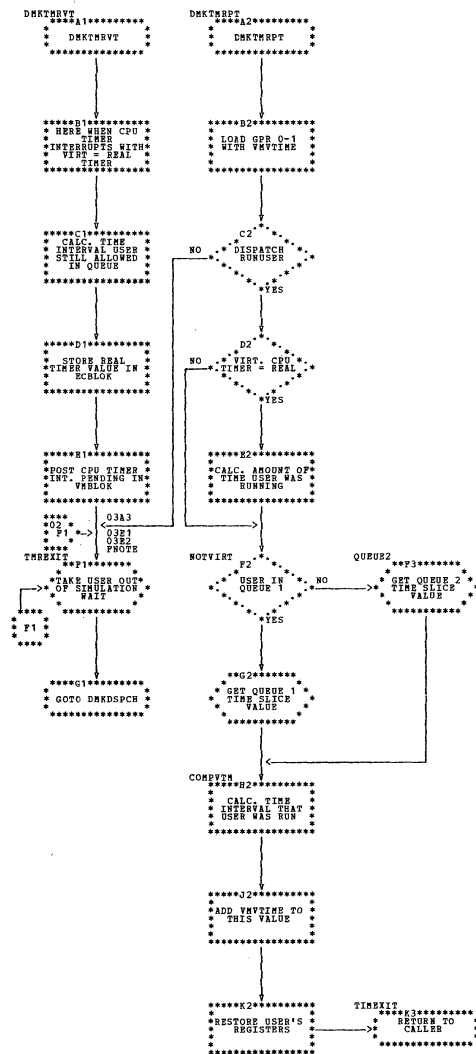
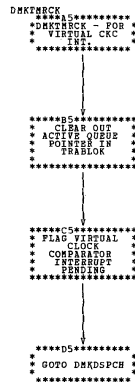
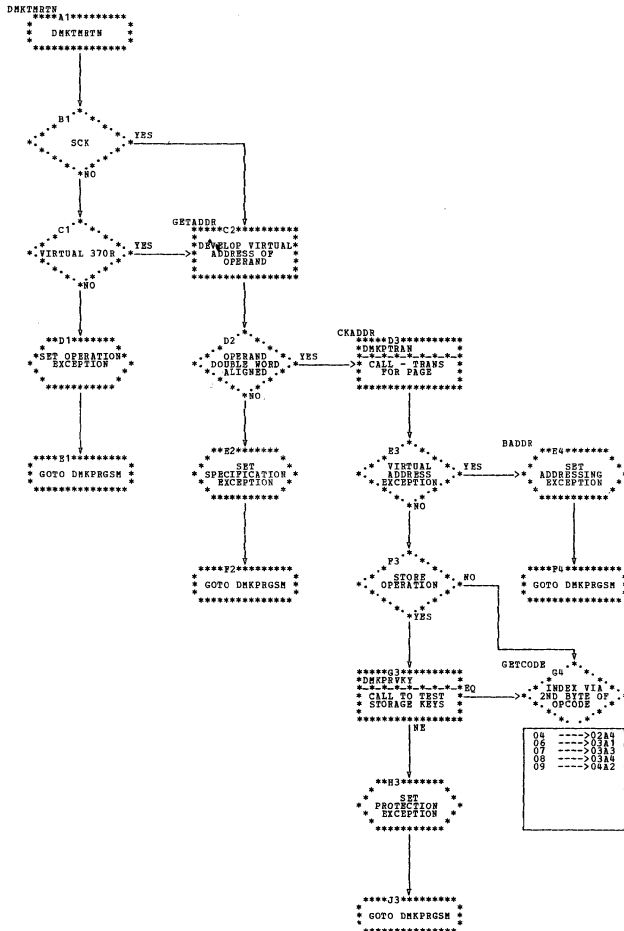
70:05  
0963







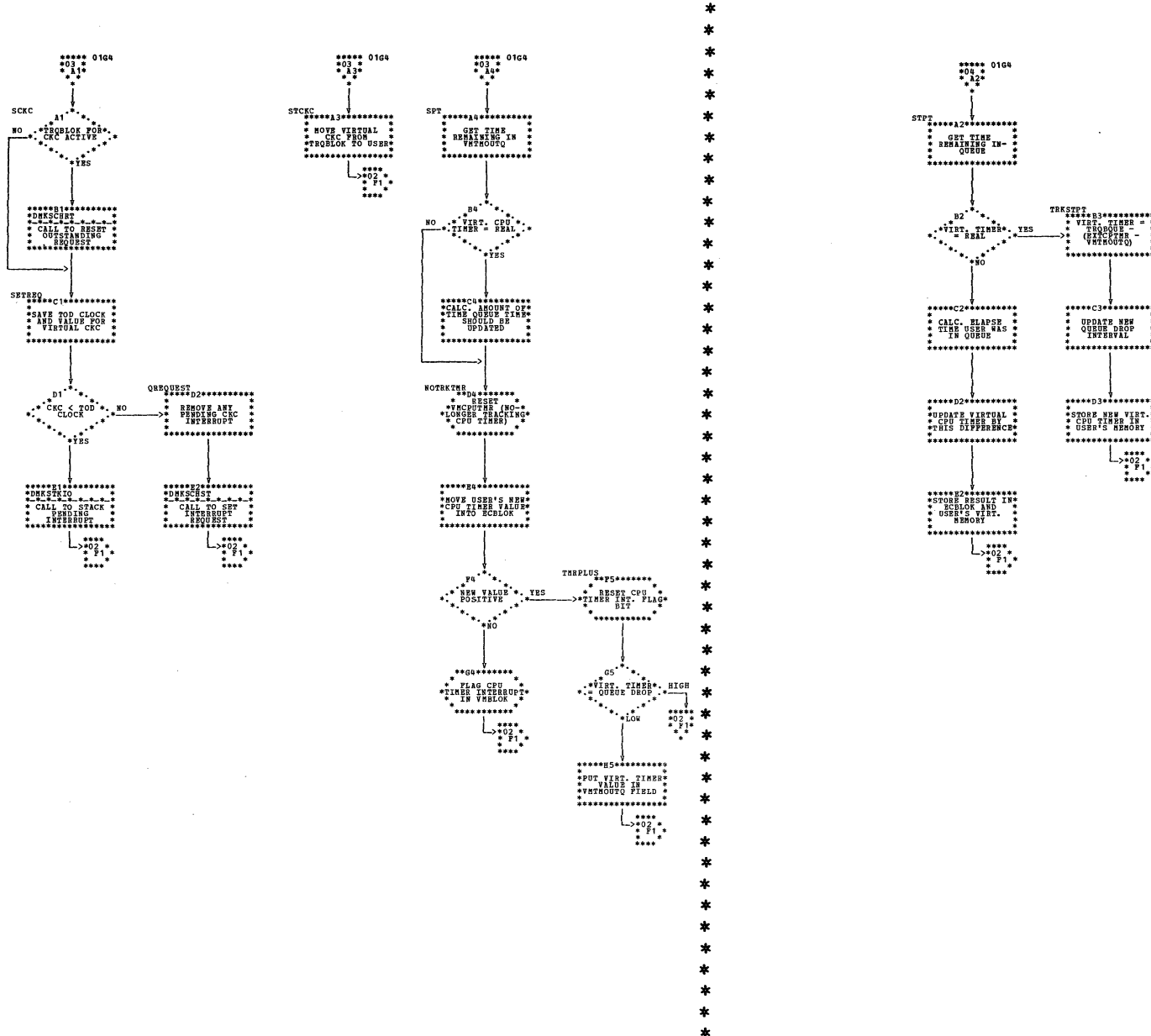




TO: F1  
C365  
C365  
C365  
C365

| DMKTRM -- Virtual CPU Timer and Clock Comparator Simulator (Parts 1 and 2 of 4)

DMKTMR -- Virtual CPU Timer and Clock Comparator Simulator (Parts 3 and 4 of 4)

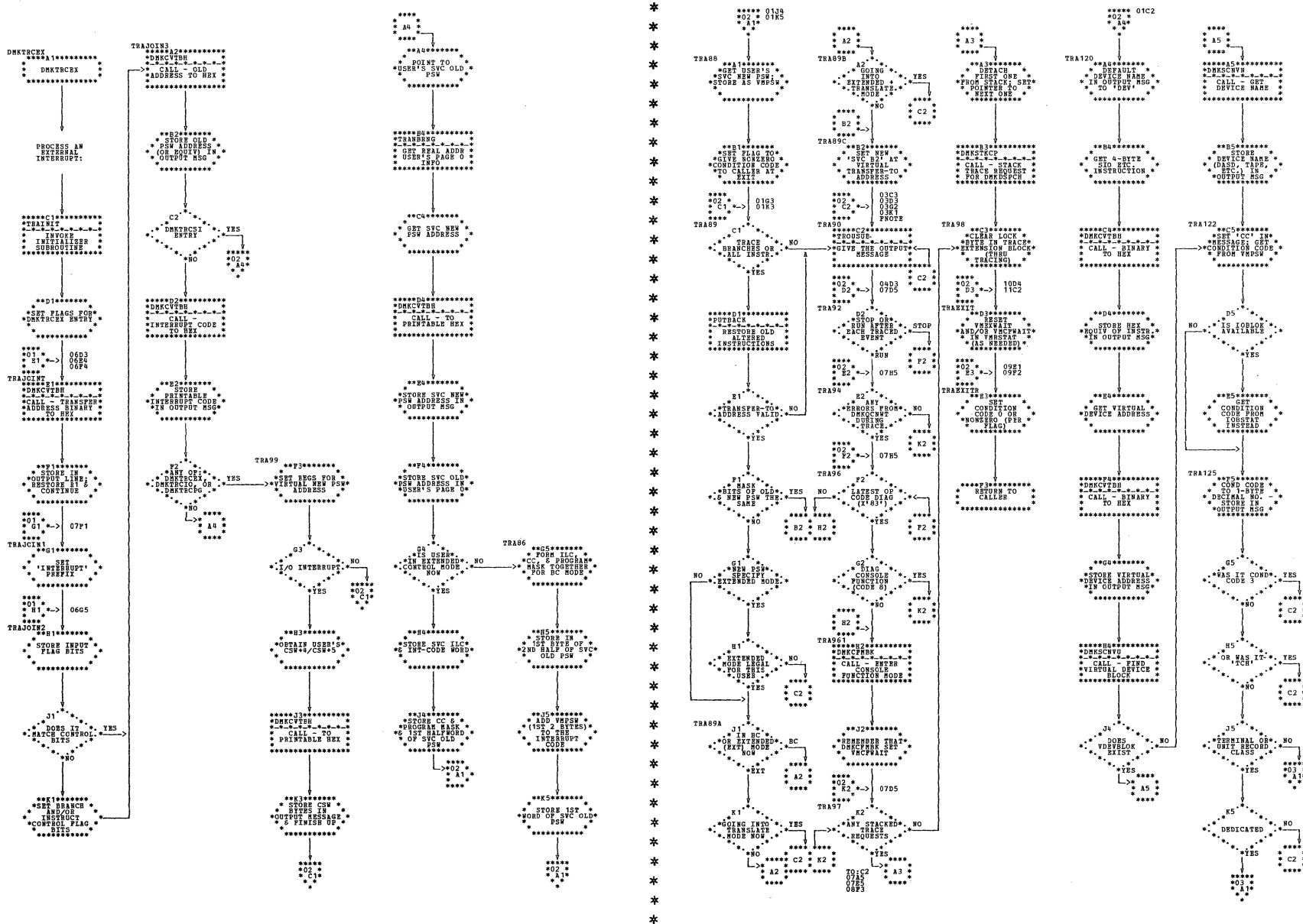








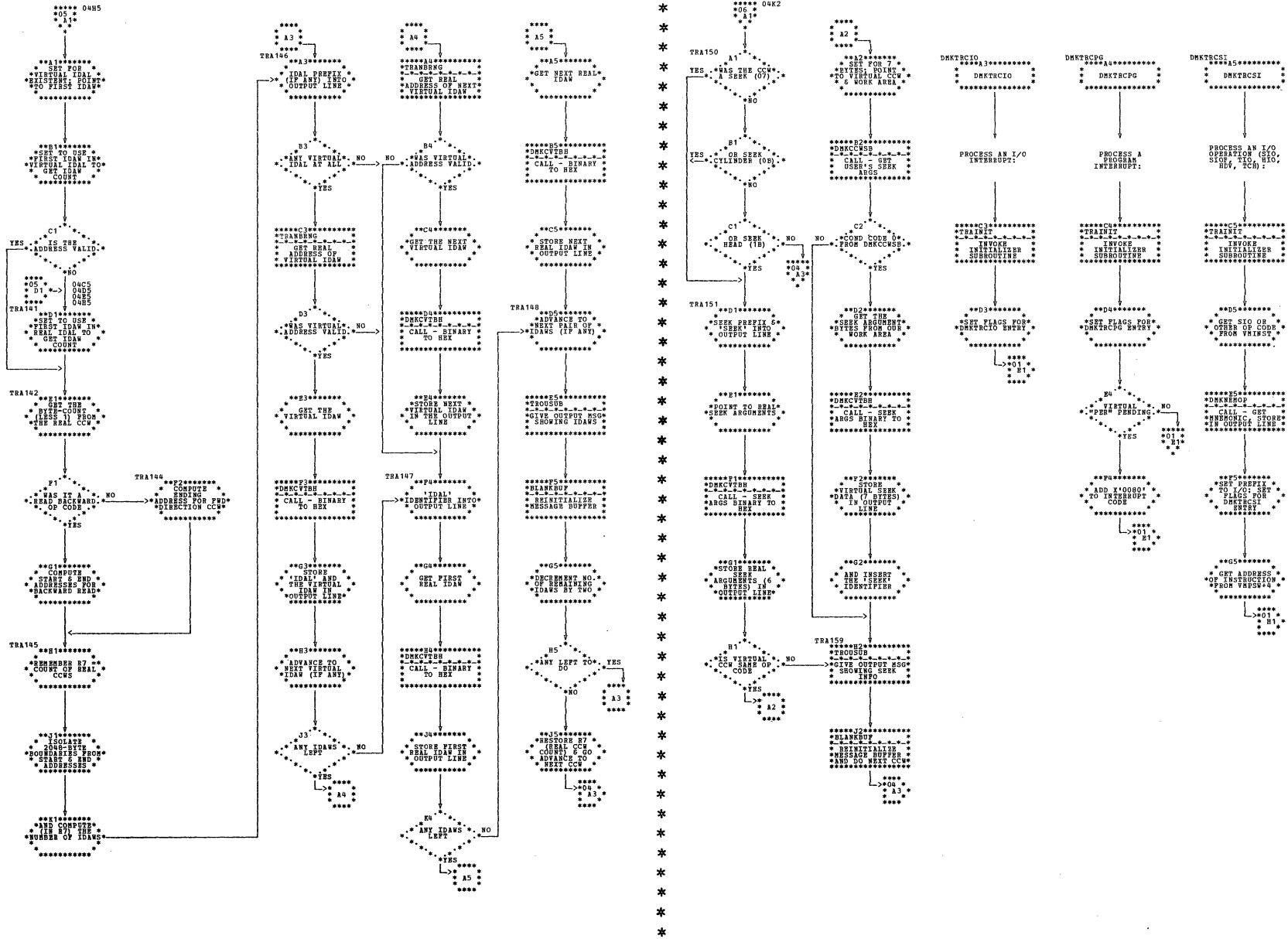
DMKTRC -- Virtual Machine Tracing Supervisor (Parts 1 and 2 of 13)

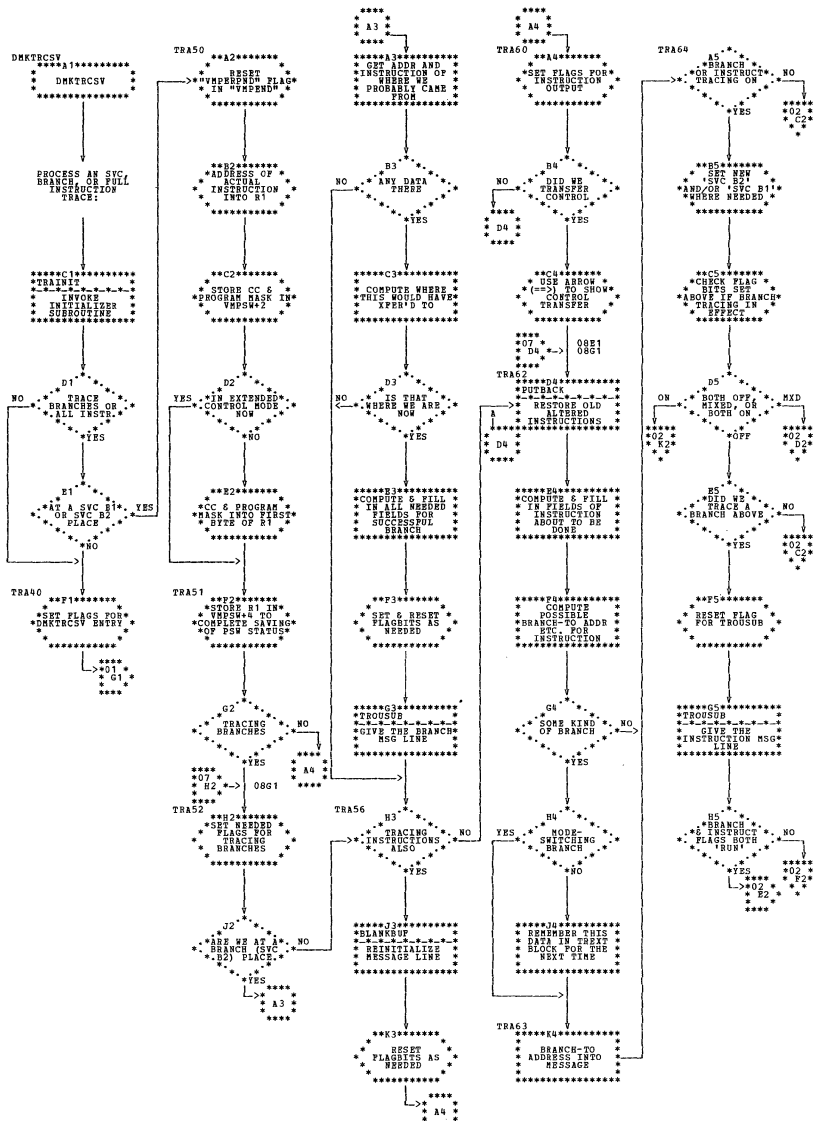




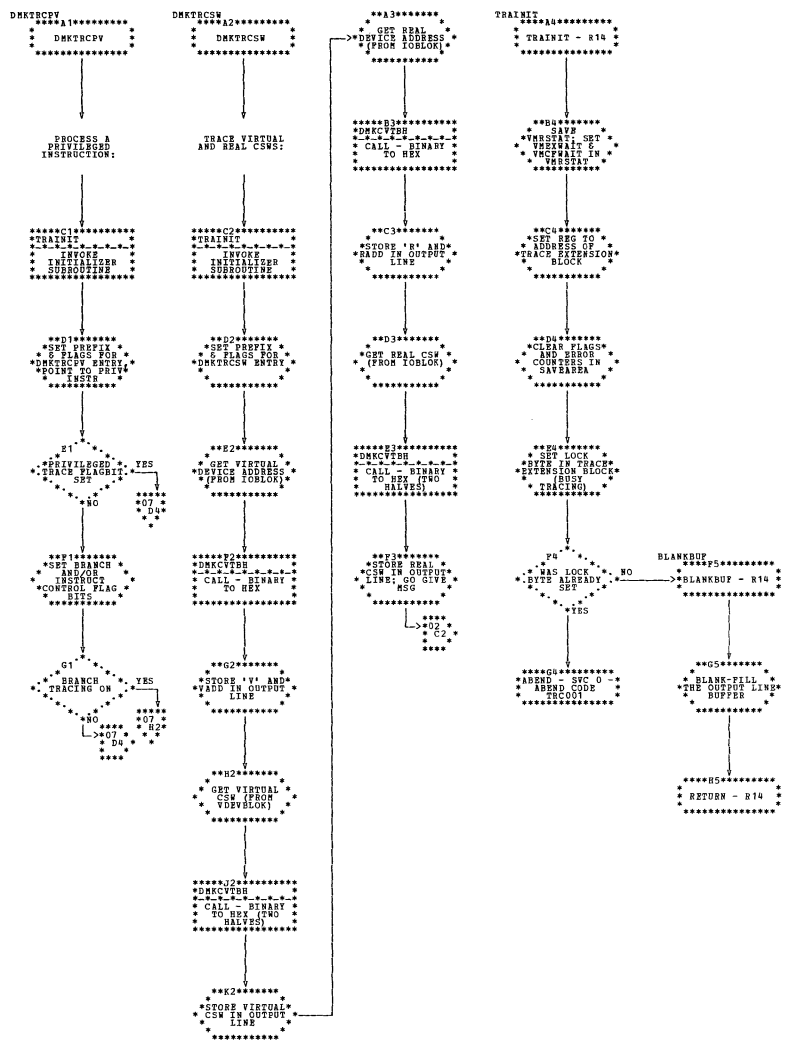


DMKTRC -- Virtual Machine Tracing Supervisor (Parts 5 and 6 of 13)



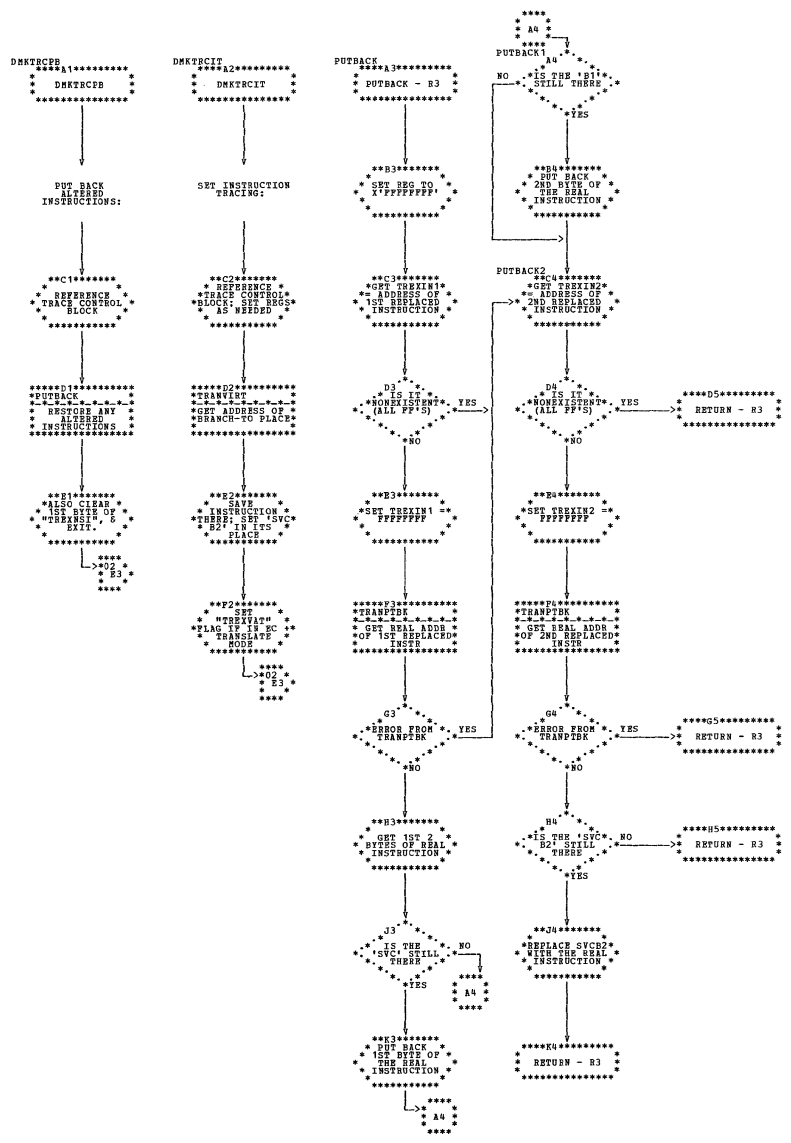


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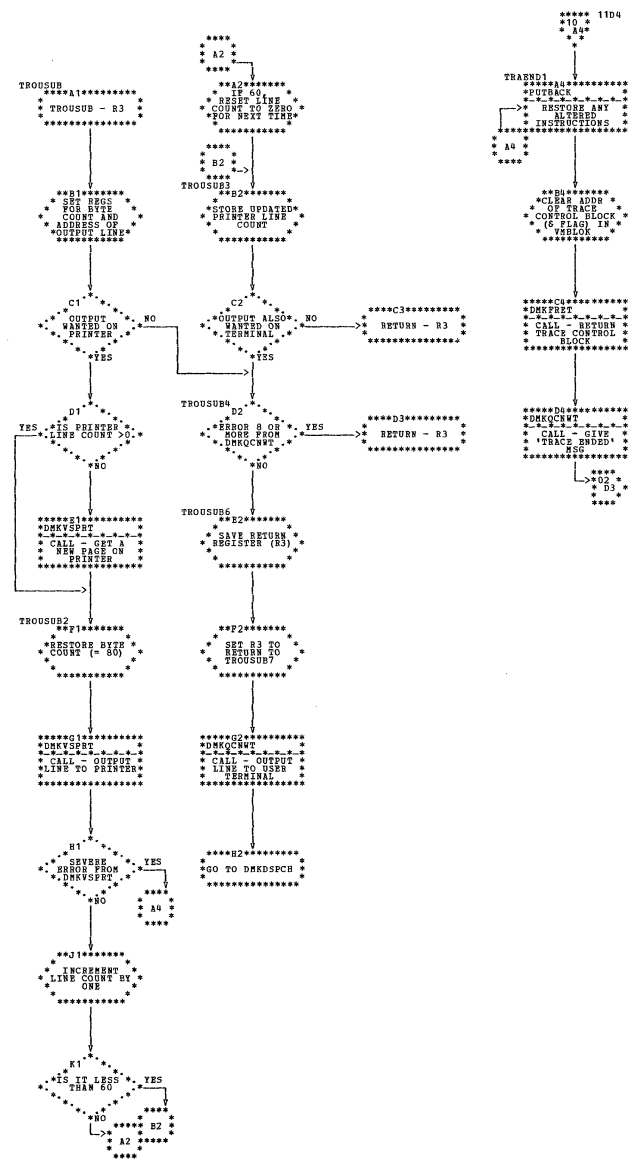


DMKTRC -- Virtual Machine Tracing Supervisor (Parts 7 and 8 of 13)

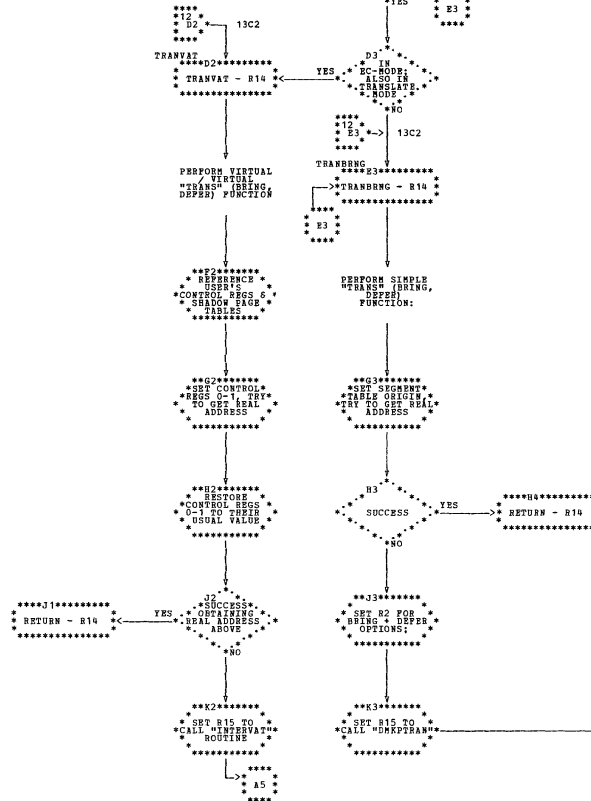
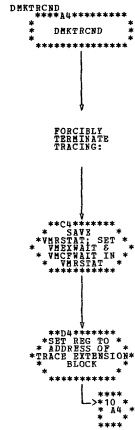
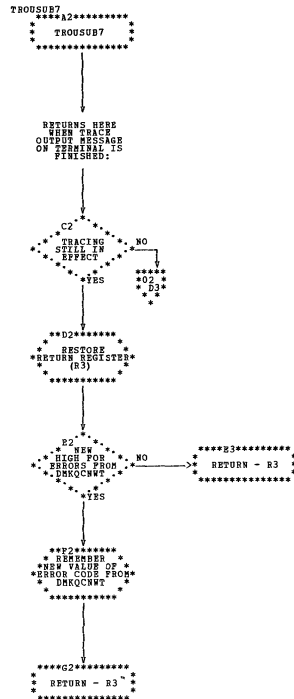
| DMKTRC -- Virtual Machine Tracing Supervisor (Parts 9 and 10 of 13)



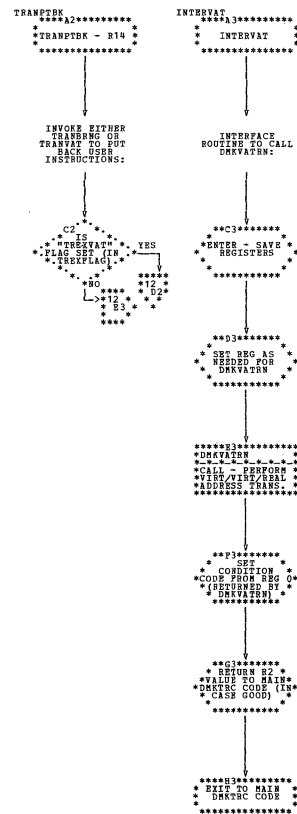
\* \* \* \* \*

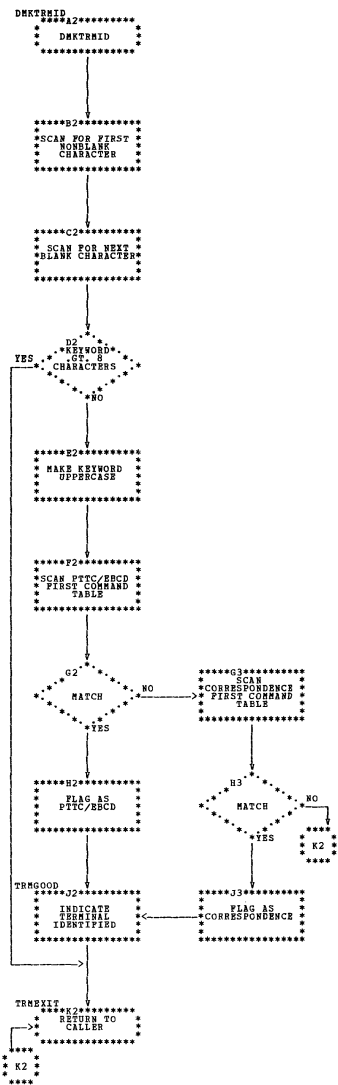


\* \* \* \* \*



| DMKTRC -- Virtual Machine Tracing Supervisor (Part 13 of 13)

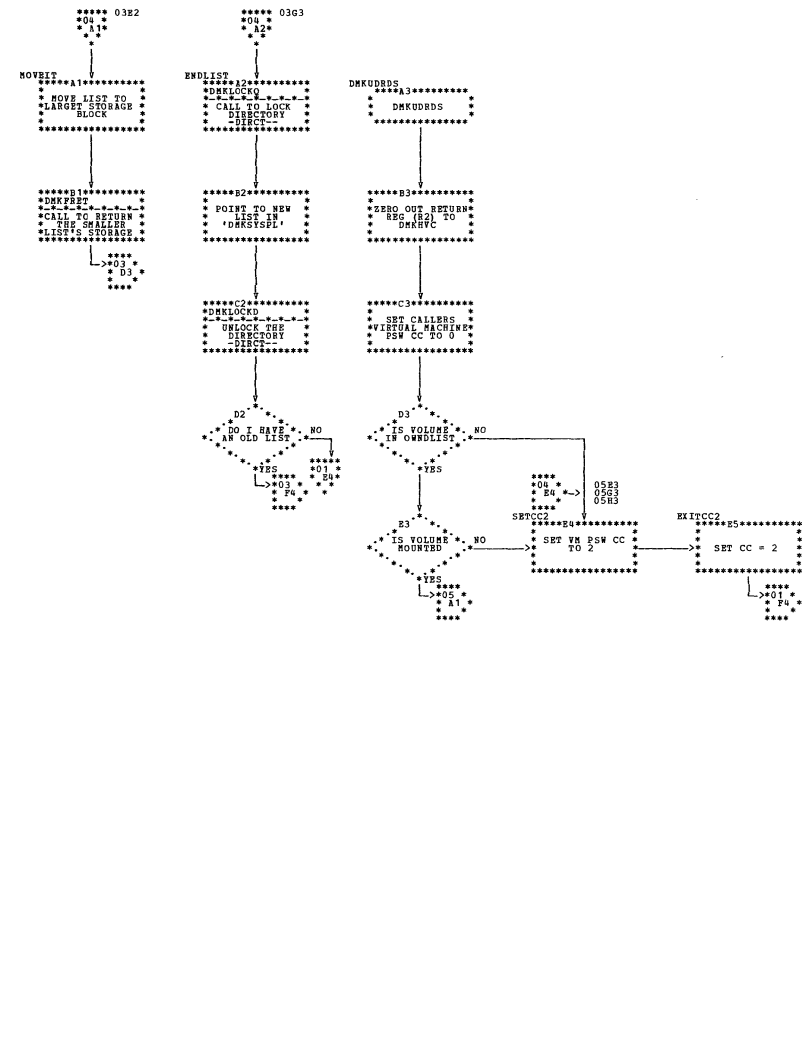
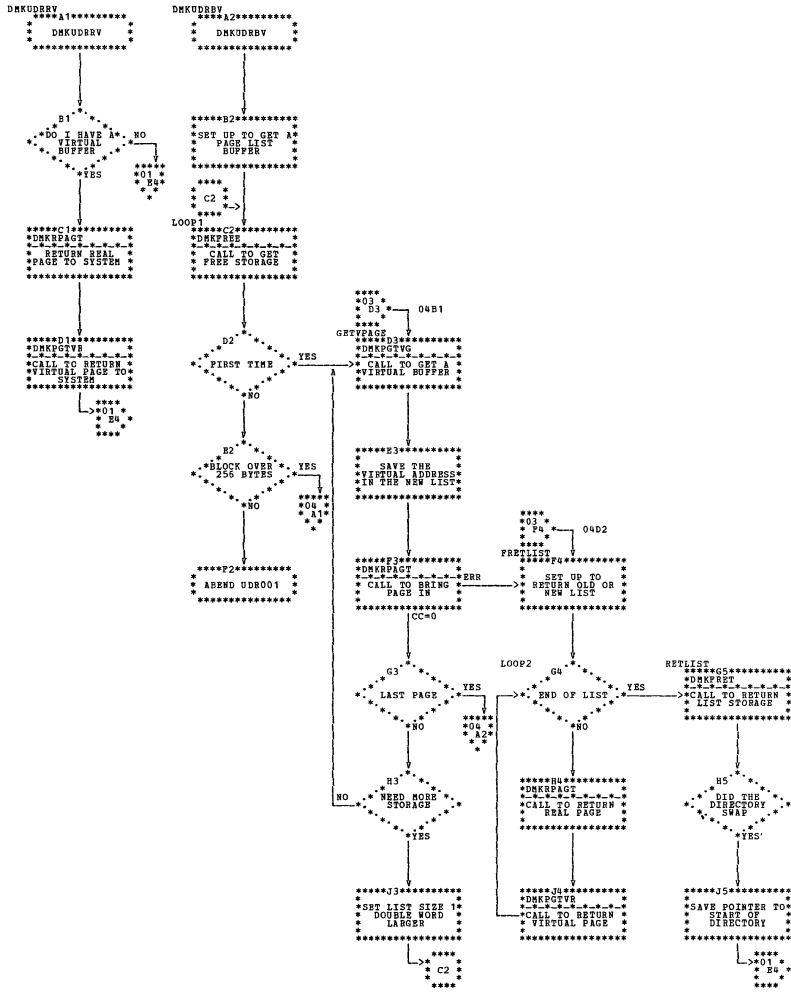




DMKTRM -- Identify Terminal (Part 1 of 1)

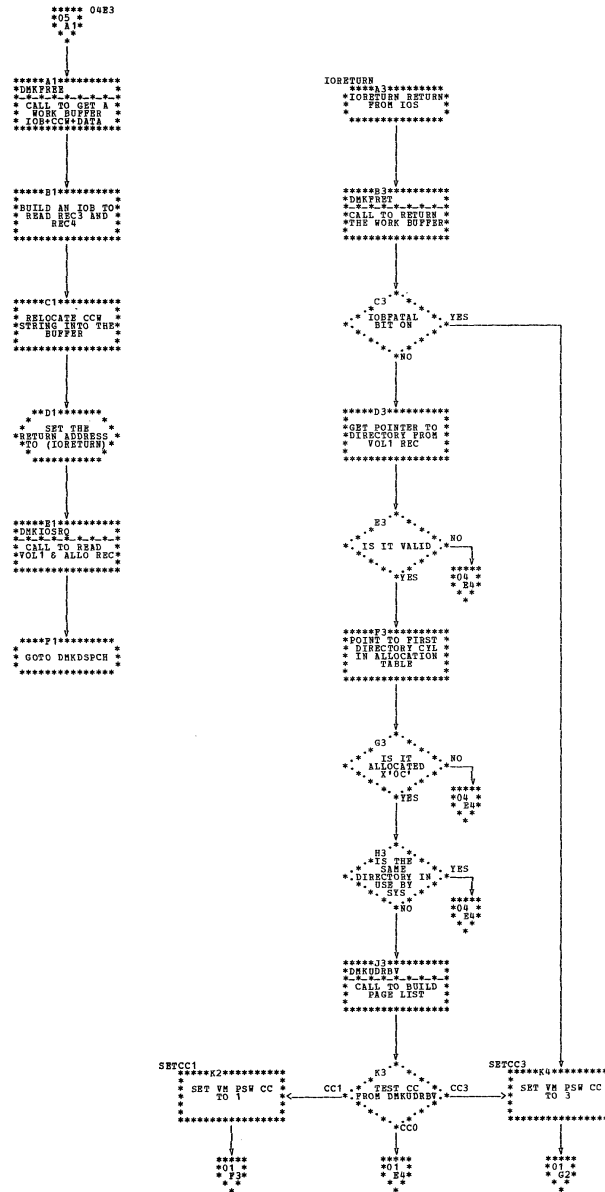


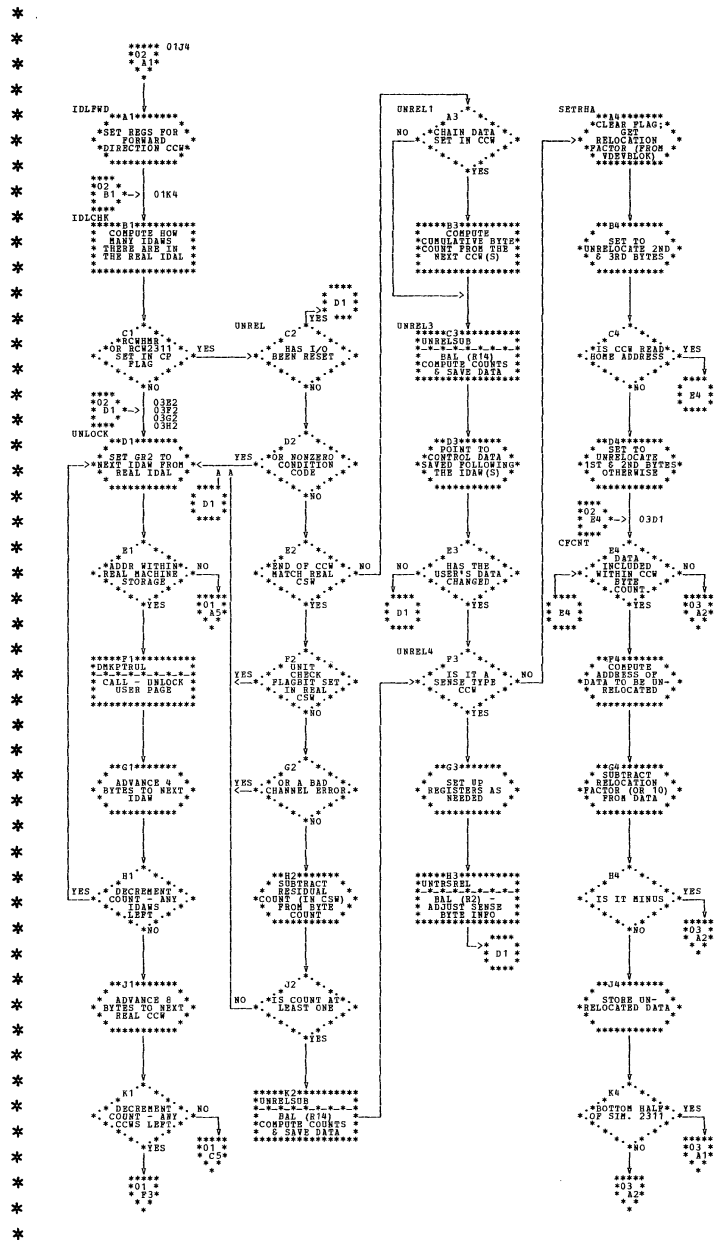
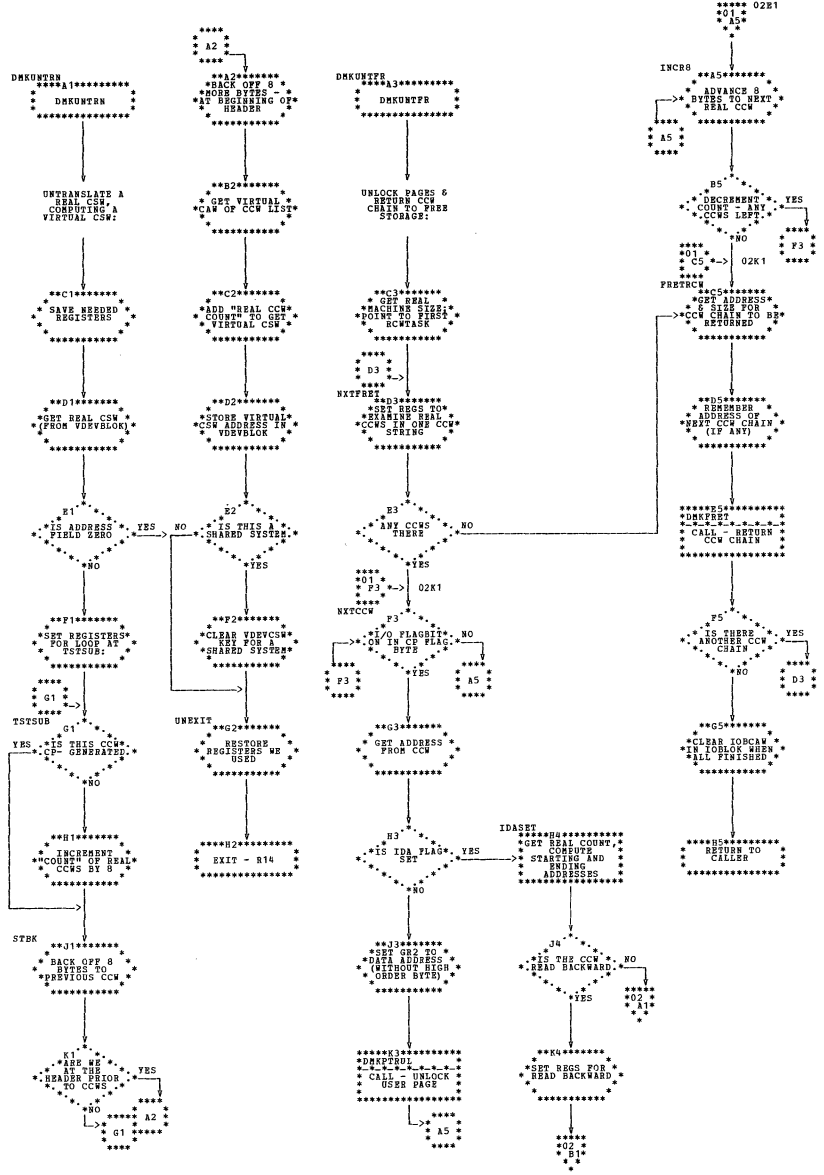




| DMKUDR -- User Directory Manager (Parts 3 and 4 of 5)

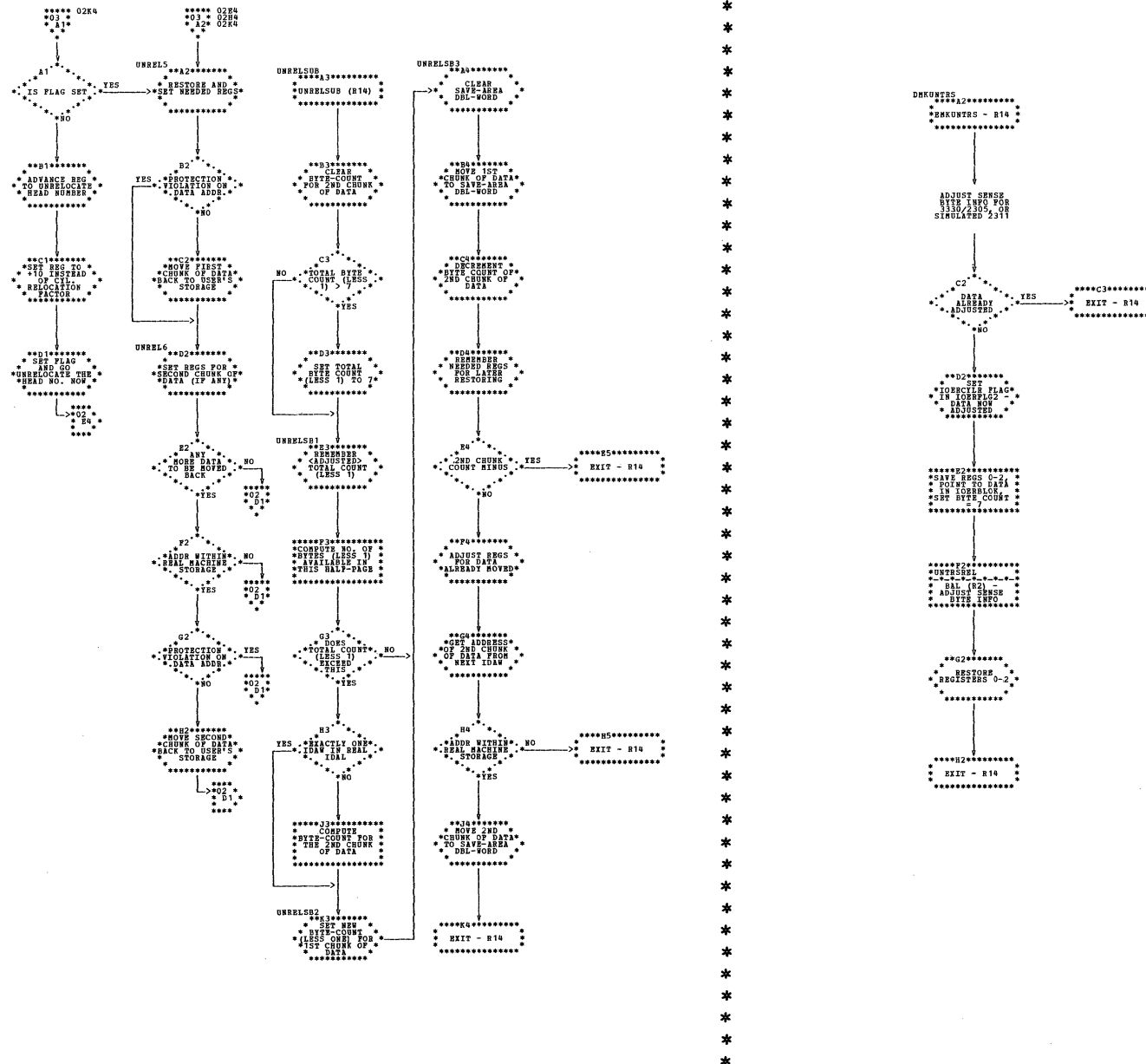
| DMKUDR -- User Directory Manager (Part 5 of 5)





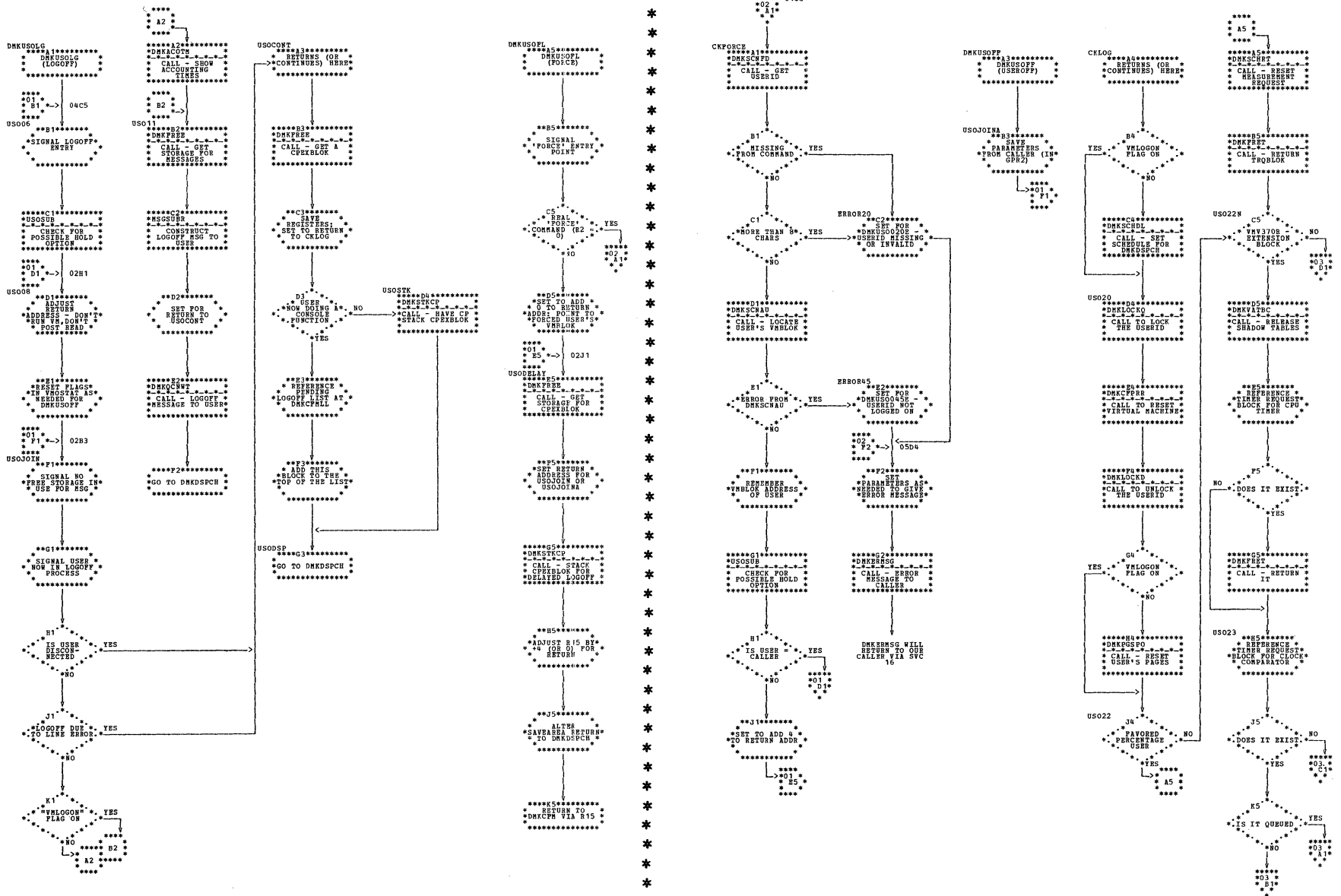
DMKUNT -- Virtual I/O Subroutines (Parts 1 and 2 of 6)

DMKUNT -- Virtual I/O Subroutines (Parts 3 and 4 of 6)



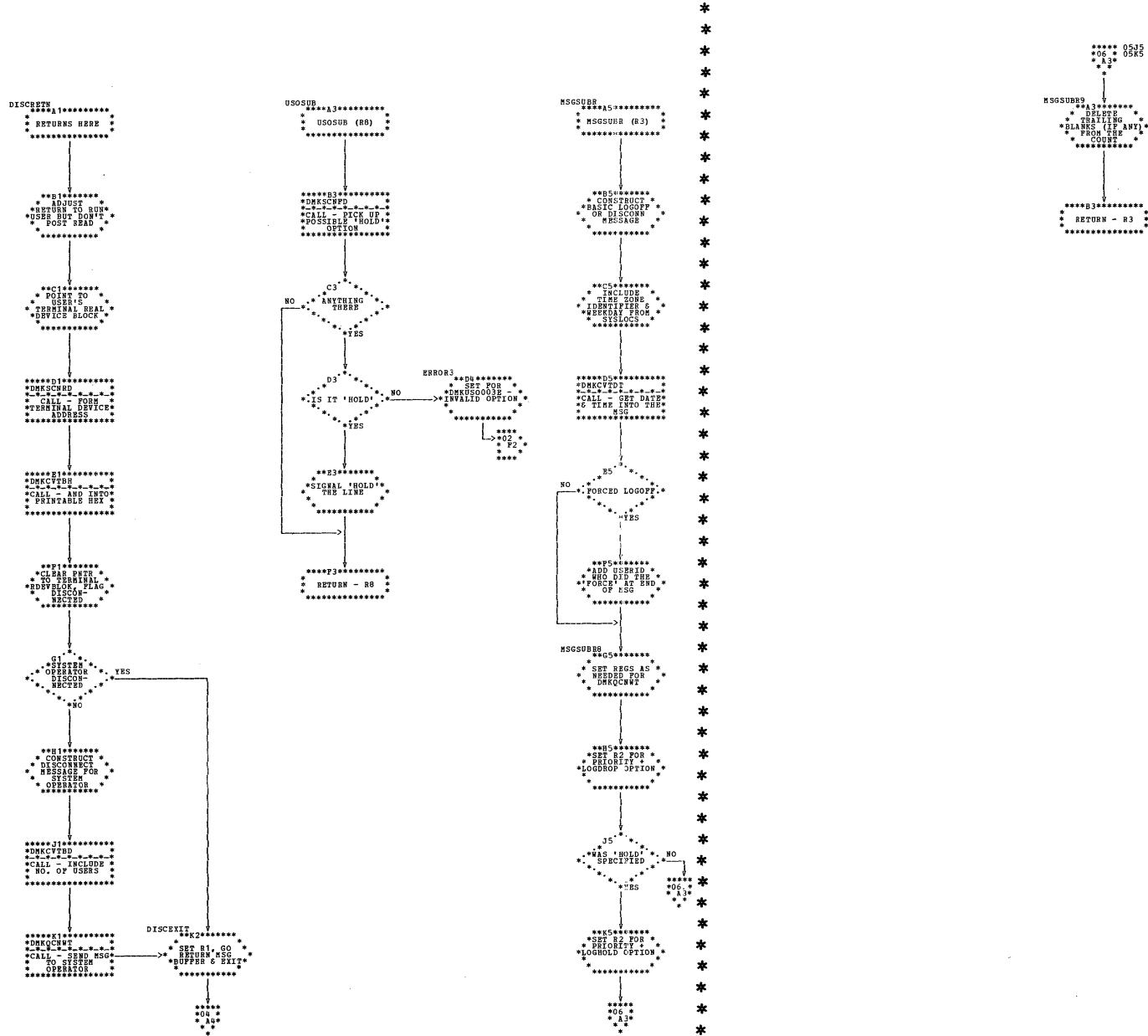


DMKUSO -- Process DISCON, FORCE, and LOGOFF Commands; Logoff Routine (Parts 1 and 2 of 6)

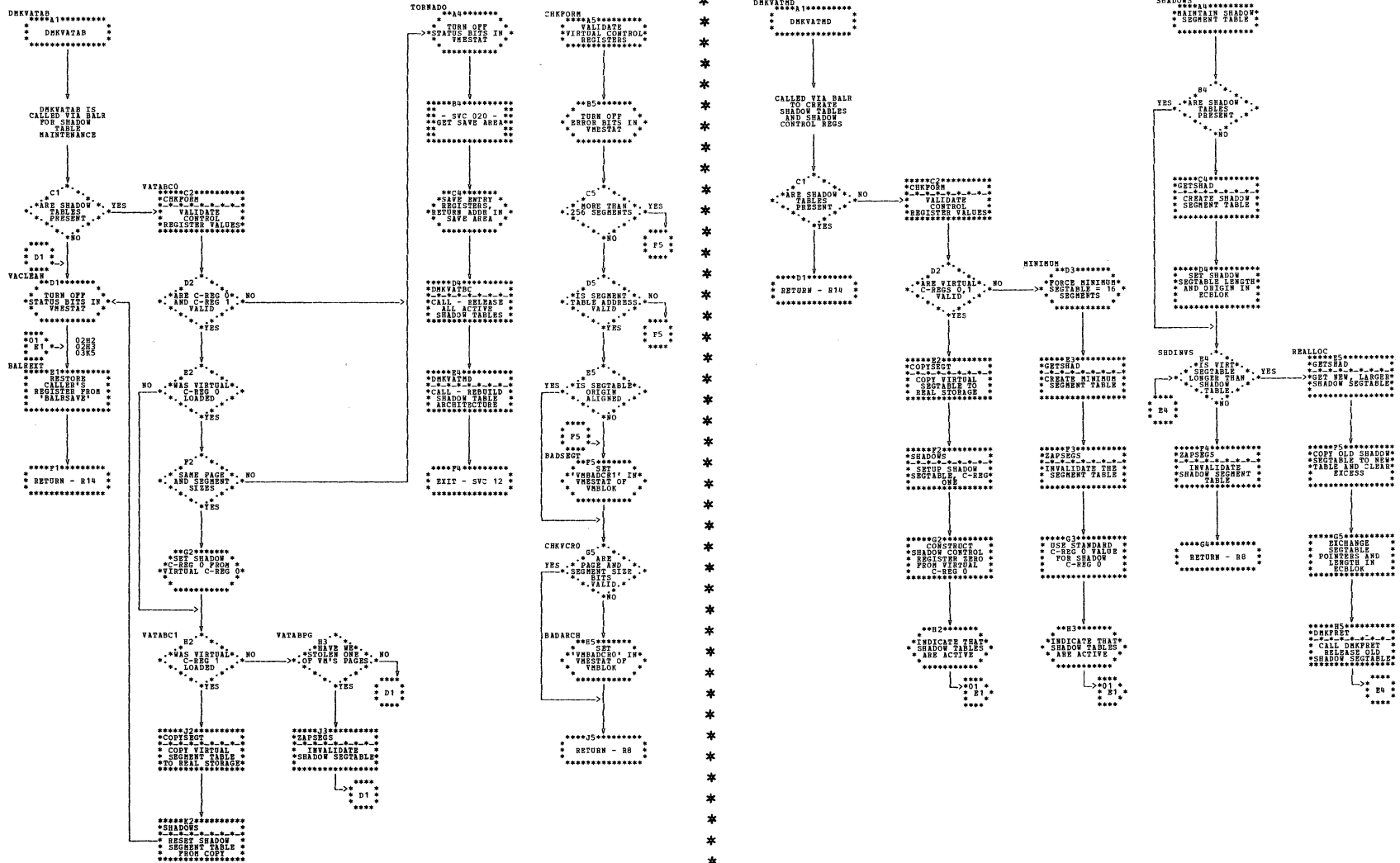




DMKUSO -- Process DISCON, FORCE, and LOGOFF Commands; Logoff Routine (Parts 5 and 6 of 6)

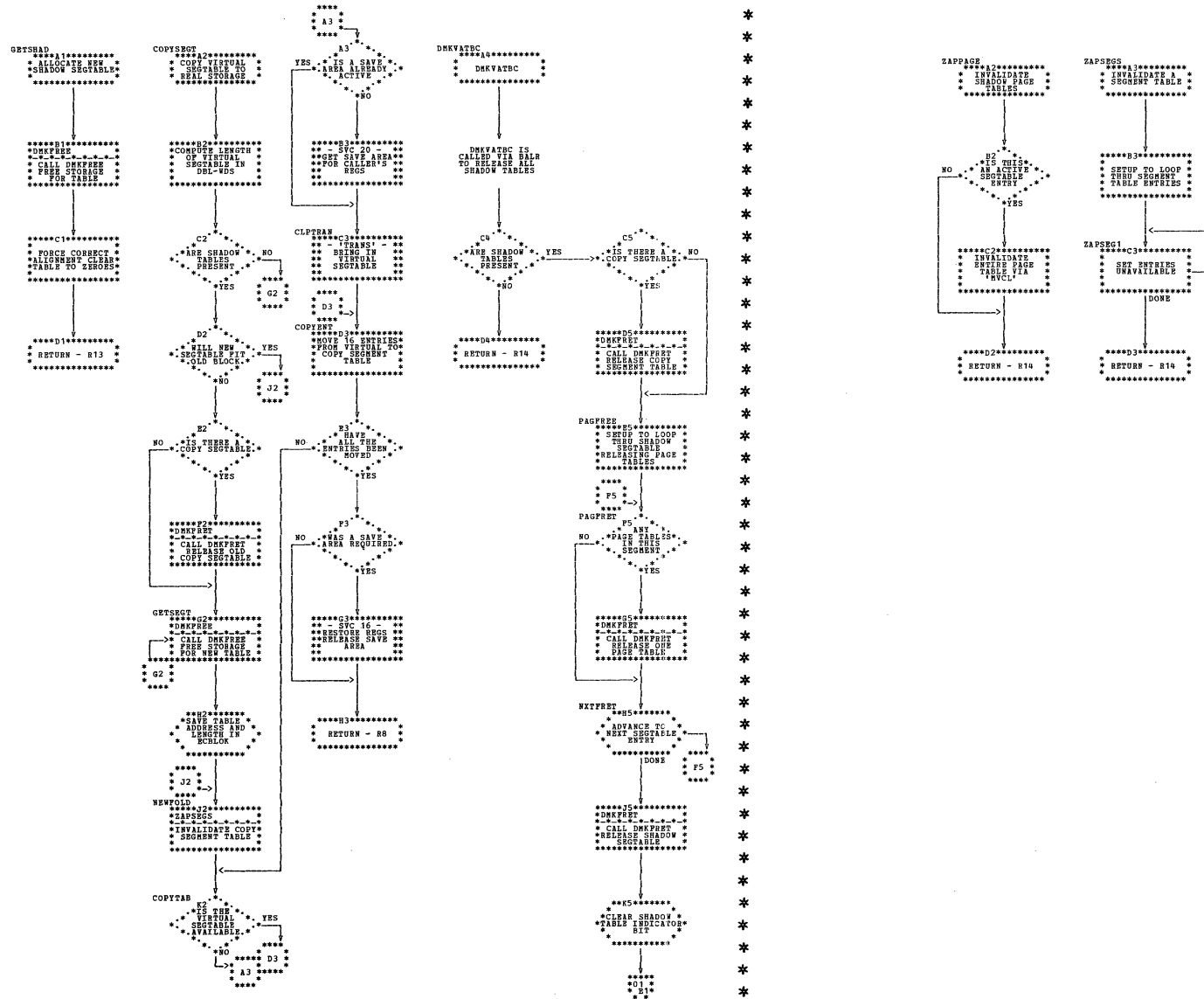






DMKVAT -- Virtual Storage Manager for EC Mode Virtual Machine that does Paging (Parts 1 and 2 of 11)

DMKVAT -- Virtual Storage Manager for EC Mode Virtual Machine that does Paging (Part 3 and 4 of 11)



```

DMKVATLA
*****A1*****
  DMKVATLA
  
```

```

DMKVATLA IS
CALLED VIA SVC
FROM DMKPRVLE
FOR I/O
SIMULATION
  
```

```

**C1**
SVC SWITCH
INDICATING
SINGLE LEVEL
TRANSLATION
  
```

```

****
D1  -> 07C4
  
```

```

REJOIN D1
  
```

```

YES -> ARE SHADOW
      PRESENT
NO
  
```

```

*****H1*****
DMKVATMD
CALL DMKVATD
ALLOCATE SHADOW
TABLES
  
```

```

VATNEXT
*****I1*****
SETUPAT
ACCESS ECBLOC
COPY SHADOWS
ARCHITECTURE
  
```

```

G1 -> ARE C-REG
      VALUES VALID
NO
  
```

```

TRNCHK G2 -> IS THIS
        TABL
        SIMULATION
  
```

```

YES
  
```

```

*****H1*****
SVCCHK
SETUP ADDRESS OF
PAGE TABLE
SET
  
```

```

G1 -> IS THE
      VIRTUAL
      PAGE TABLE
      AVAILABLE
YES
NO
  
```

```

*****
D1  -> 07C7
  
```

```

*****H2*****
SVC 16
DON'T RETURN
TO CALLER
  
```

```

****
D3  -> 09D1
      09D4
  
```

```

TRNCHK G3 -> SET
          INTERRUPT
          CODE IN
          TRANSLATION
          ERROR
  
```

```

****
D4  -> 09C2
      09C3
      09C4
      10B2
  
```

```

VATZAP G4 -> - 'TRANS' -
          SET
          FOR VIRTUAL
          PAGE ZERO
  
```

```

G4 -> TRANSLATION
        EXCEPTION
        ADDRESS IN
        PAGE ZERO
  
```

```

****K4*****
GOTO DMKPRGSE
SIMULATE ERROR
INTERRUPT
  
```

\*\*\*\*\*

```

**** 05J1
D6
A1
  
```

```

*****I1*****
- 'TRANS' -
SET ERROR
RETURN FOR
CHECKER TO
TRNCHK
  
```

```

GIVEB01
GIVEINT
  
```

```

*****C1*****
CHECKPT
IS THE VIRTUAL
PAGE AVAILABLE
  
```

```

NO -> CODE IN
      PAGE
      EXCEPTION
  
```

```

YES
  
```

```

D1 -> IS THIS
      TABL
      SIMULATION
  
```

```

YES -> COMPUTE VIRTUAL
        ADDRESS AND
        RETURN RESULTS
  
```

```

NO
  
```

```

*****I1*****
TRANSLATE
VIRTUAL
ADDRESS AND
PAGE ZERO
  
```

```

*****I1*****
SET RESULT IN
SAVE AREA FOR
CALLER
  
```

```

*****G1*****
SHADSET
UPDATE SHADOW
PAGE TABLE
  
```

```

TRNEXIT
EXIT - SVC 12
  
```

```

****
D2  -> 07A3
  
```

```

*****C2*****
CHECKPT
IS THE VIRTUAL
PAGE AVAILABLE
  
```

```

NO -> CODE IN
      PAGE
      EXCEPTION
  
```

```

YES
  
```

```

D1 -> IS THIS
      TABL
      SIMULATION
  
```

```

YES -> COMPUTE VIRTUAL
        ADDRESS AND
        RETURN RESULTS
  
```

```

NO
  
```

```

*****I1*****
TRANSLATE
VIRTUAL
ADDRESS AND
PAGE ZERO
  
```

```

*****I1*****
SET RESULT IN
SAVE AREA FOR
CALLER
  
```

```

*****G1*****
SHADSET
UPDATE SHADOW
PAGE TABLE
  
```

```

****
D3  -> 07C3
  
```

```

*****C3*****
CHECKPT
IS THE VIRTUAL
PAGE AVAILABLE
  
```

```

NO -> CODE IN
      PAGE
      EXCEPTION
  
```

```

YES
  
```

```

D1 -> IS THIS
      TABL
      SIMULATION
  
```

```

YES -> COMPUTE VIRTUAL
        ADDRESS AND
        RETURN RESULTS
  
```

```

NO
  
```

```

*****I1*****
TRANSLATE
VIRTUAL
ADDRESS AND
PAGE ZERO
  
```

```

*****I1*****
SET RESULT IN
SAVE AREA FOR
CALLER
  
```

```

*****G1*****
SHADSET
UPDATE SHADOW
PAGE TABLE
  
```

```

**** 05G2
D6
D7
D8
  
```

```

*****I1*****
SET
INTERRUPT
CODE IN
TRANSLATION
ERROR
  
```

```

GIVEB02
GIVEINT
  
```

```

*****C1*****
CHECKPT
IS THE VIRTUAL
PAGE AVAILABLE
  
```

```

NO -> CODE IN
      PAGE
      EXCEPTION
  
```

```

YES
  
```

```

D1 -> IS THIS
      TABL
      SIMULATION
  
```

```

YES -> COMPUTE VIRTUAL
        ADDRESS AND
        RETURN RESULTS
  
```

```

NO
  
```

```

*****I1*****
TRANSLATE
VIRTUAL
ADDRESS AND
PAGE ZERO
  
```

```

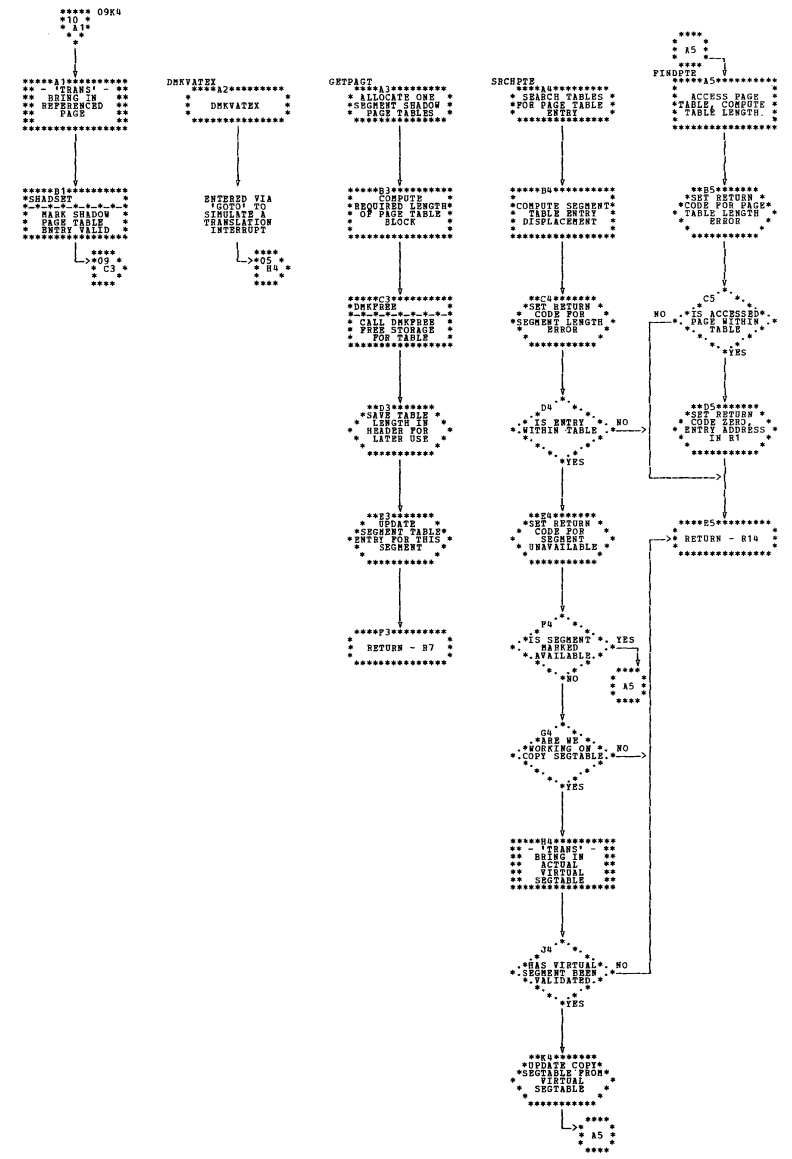
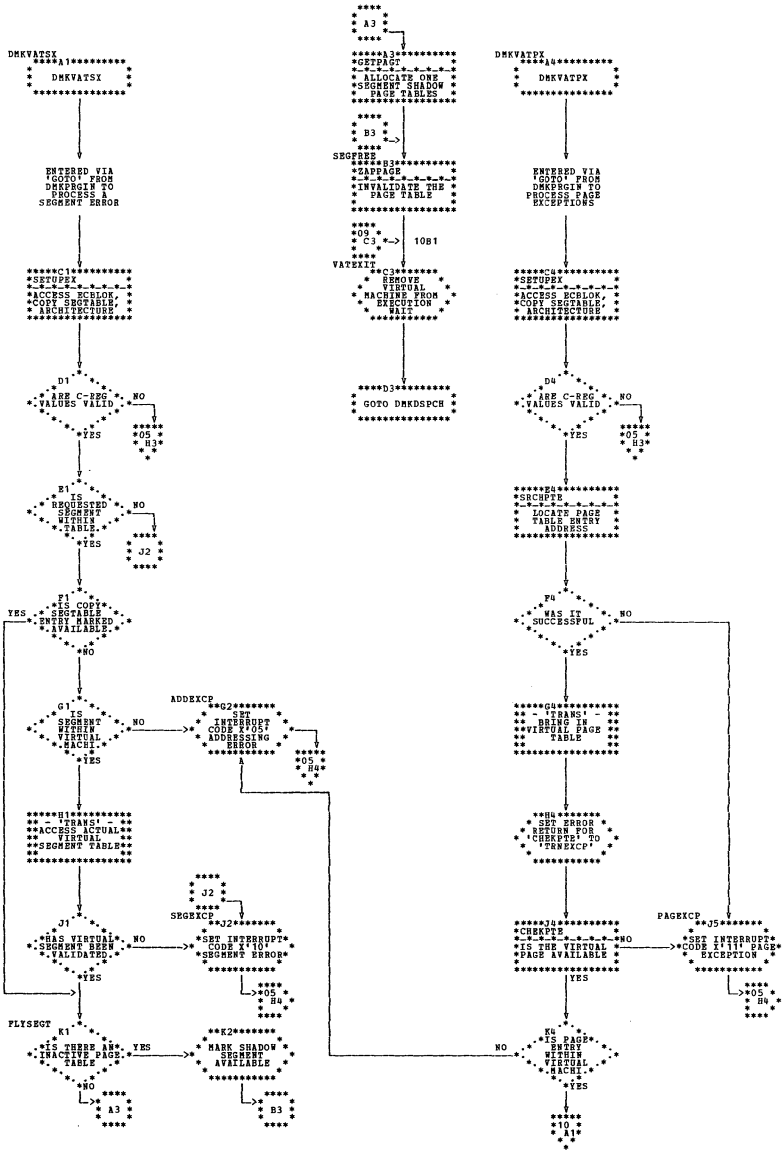
*****I1*****
SET RESULT IN
SAVE AREA FOR
CALLER
  
```

```

*****G1*****
SHADSET
UPDATE SHADOW
PAGE TABLE
  
```

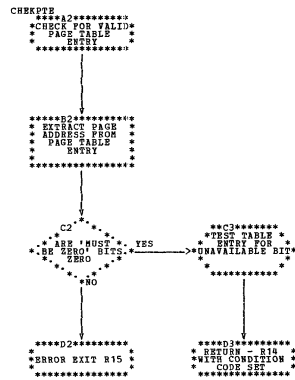
| DMKVAT -- Virtual Storage Manager for EC Mode Virtual Machine that does Paging (Part 5 and 6 of 11)





| DMKVAT -- Virtual Storage Manager for EC Mode Virtual Machine that does Paging (Parts 9 and 10 of 11)

| DMKVAT -- Virtual Storage Manager for EC Mode Virtual Machine that does Paging (Part 11 of 11)



DMKVCA  
\*\*\*\*\*  
DMKVCA

ENTERED VIA SVC  
FROM INITIALIZER  
FOR 'SIC' CCA  
VIRTUAL CCA

\*\*\*\*\*C1\*\*\*\*\*  
\*CLEAR FLAG BYT\*  
\*SERUP LABEL\*  
\*HOUSE LABEL\*  
\*ADDRESS FIELDS\*

D1  
IS THE  
VIRTUAL CCA  
COUPLED

VCASART  
\*\*\*\*\*  
\*CCTCCA\*  
\*ACCESS\*  
\*CHNLR\* AND  
\*CHNLR\*

\*\*\*\*\*F1\*\*\*\*\*  
\*IMMEDIATE\*  
\*CALL - SIGNAL\*  
\*ENDING STATUS\*  
\*NOT READY DET\*

\*\*\*\*\*F2\*\*\*\*\*  
\*SET STATUS\*  
\*ZERO SENSE\*  
\*FOR FIRST CCA\*  
\*DECODE\*

\*\*\*\*\*F3\*\*\*\*\*  
\*SAVE ADDRESS\*  
\*CHNLR\*

VCASRT  
\*\*\*\*\*  
\*EXIT - SVC 12\*

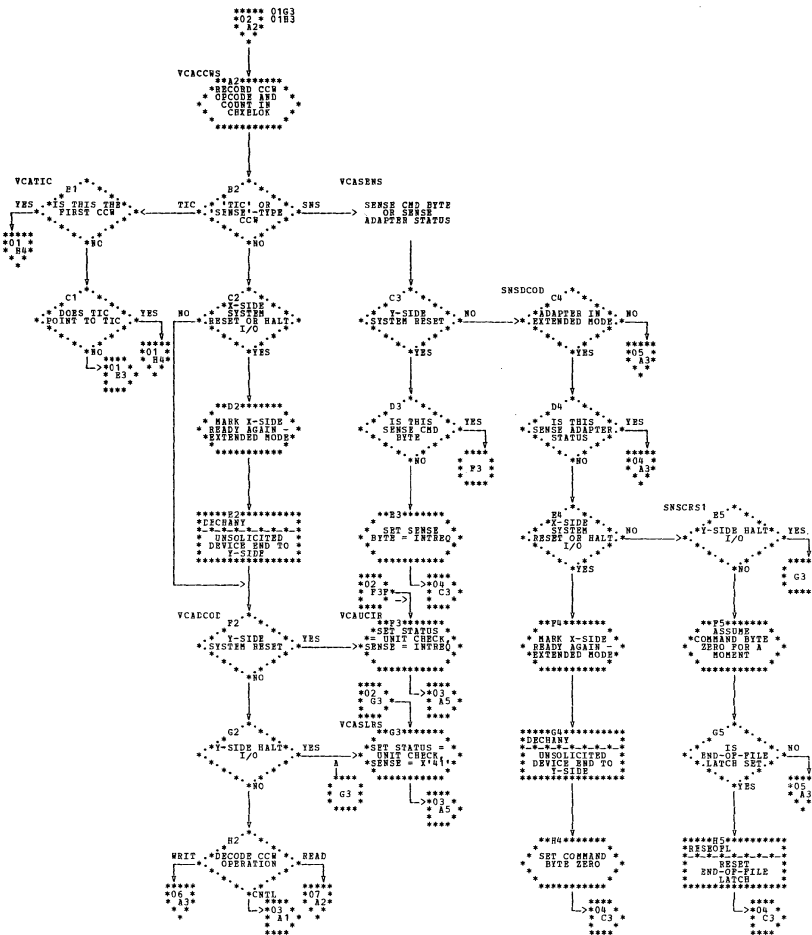
\*\*\*\*\*F4\*\*\*\*\*  
\*CHNLR\*  
\*VALIDATE REAL\*  
\*CCW\*

G3  
IS THIS THE  
FIRST CCA

H3  
ARE CAN  
BYTES 4-7 -  
ZERO

C0F1  
C0C4  
C096  
C093

\*\*\*\*\*



TO: P3 TO: G3  
1184 1183



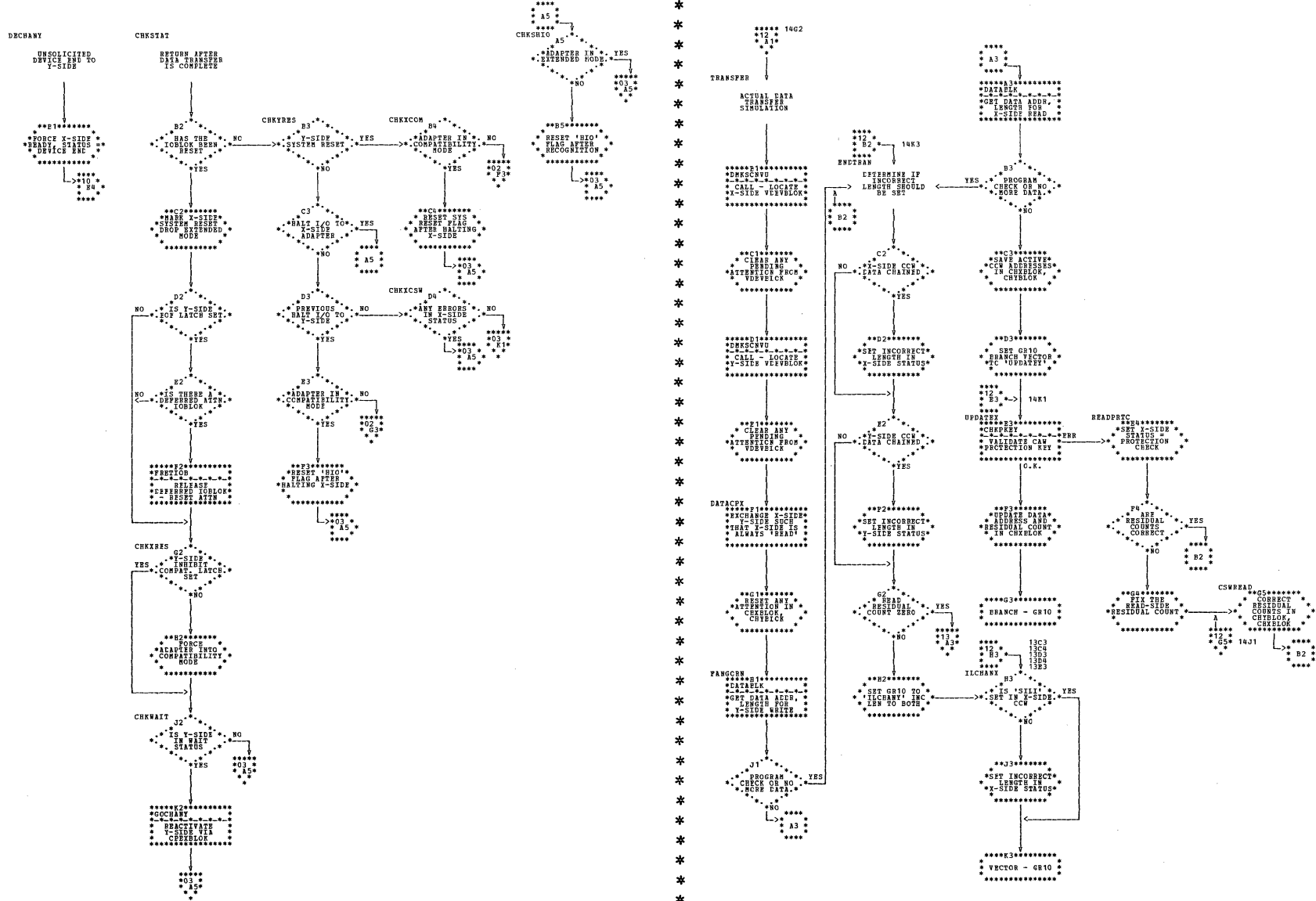


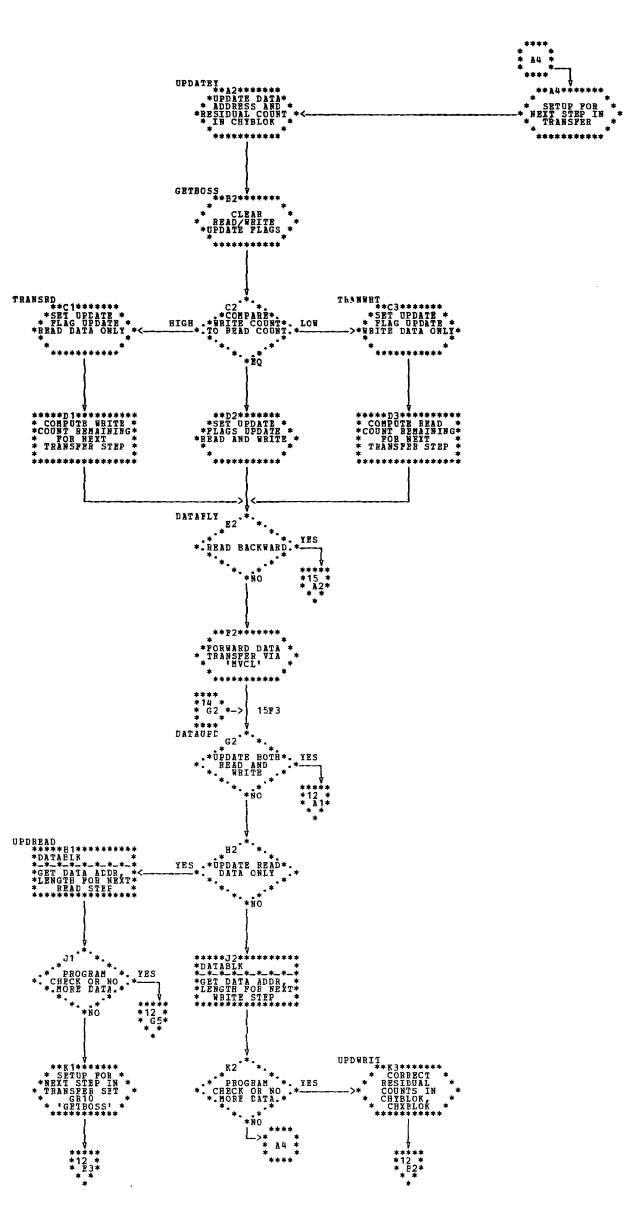
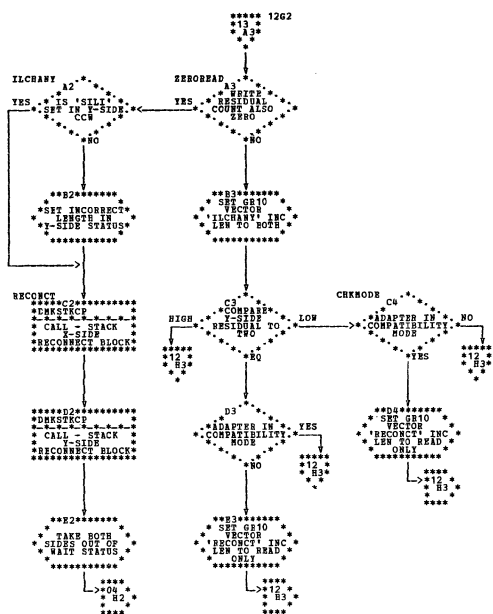






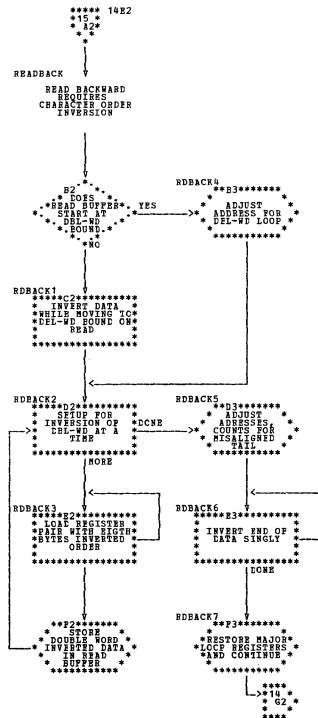
DMKVCA -- Simulate Channel-to-Channel Adapter between two Virtual Machines (Parts 11 and 12 of 19)



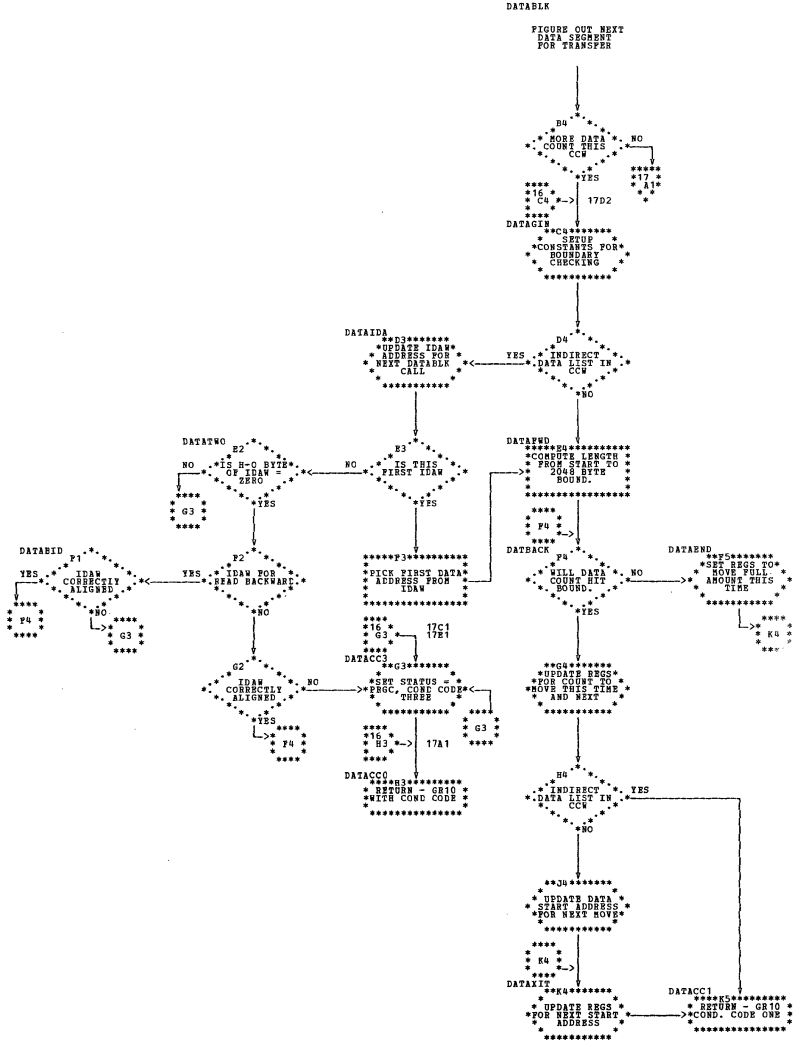


| DMKVCA -- Simulate Channel-to-Channel Adapter between two Virtual Machines (Parts 13 and 14 of 19)

DMKVCA -- Simulate Channel-to-Channel Adapter between two Virtual Machines (Parts 15 and 16 of 19)

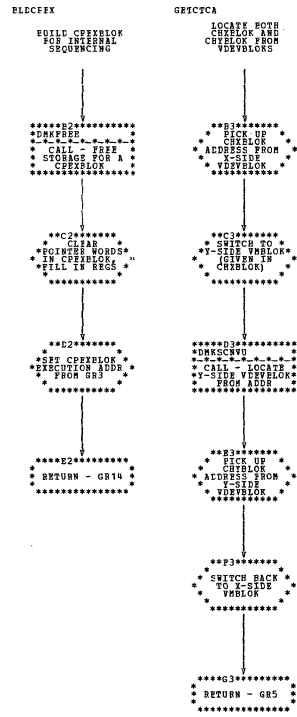


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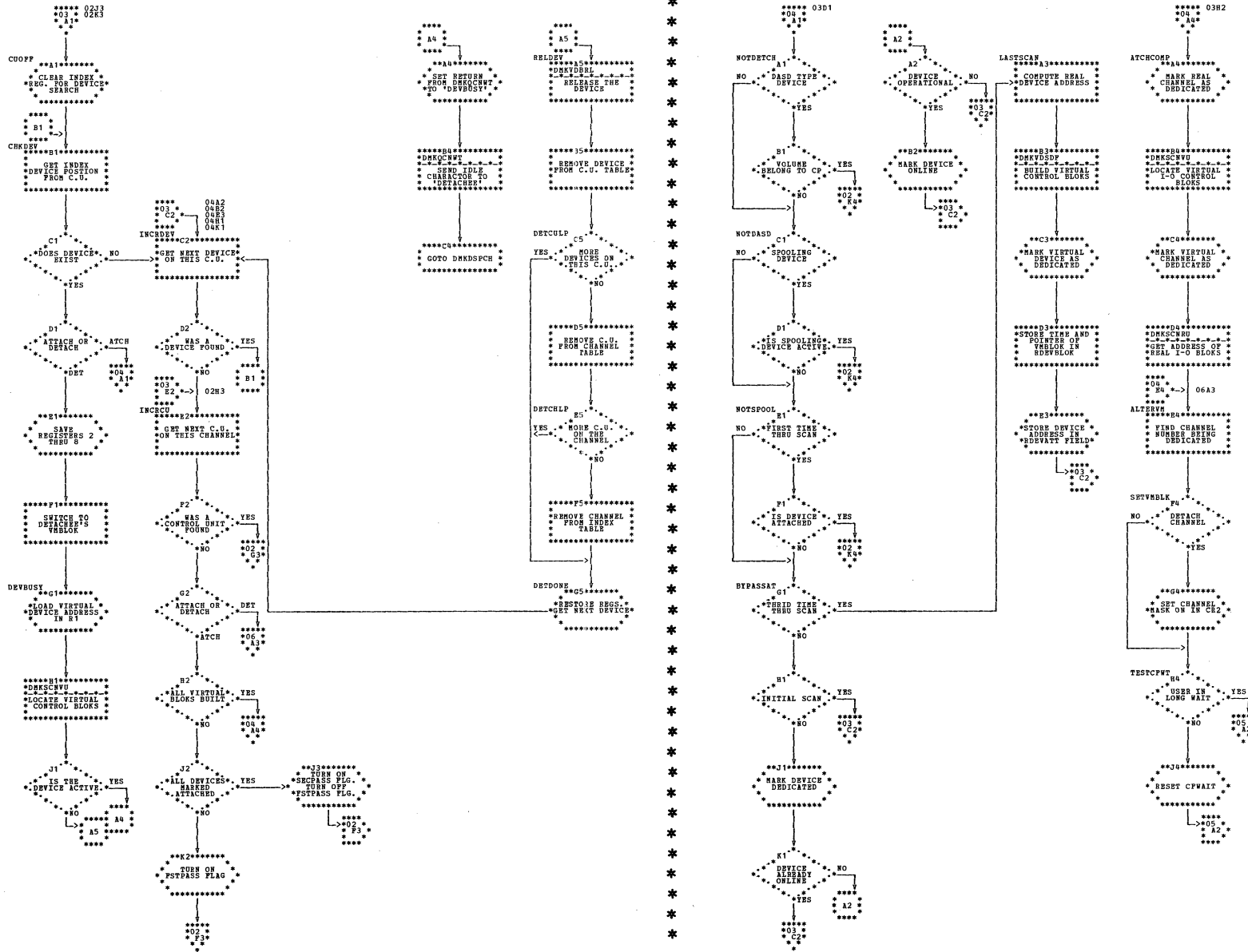
DMKVCA -- Simulate Channel-to-Channel Adapter between two Virtual Machines (Part 19 of 19)





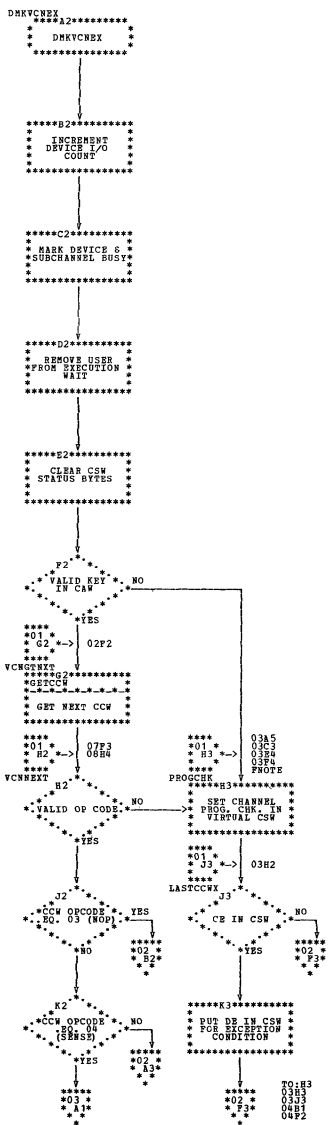


DMKVCH -- Process ATTACH and DETACH Channel Commands (Parts 3 and 4 of 6)

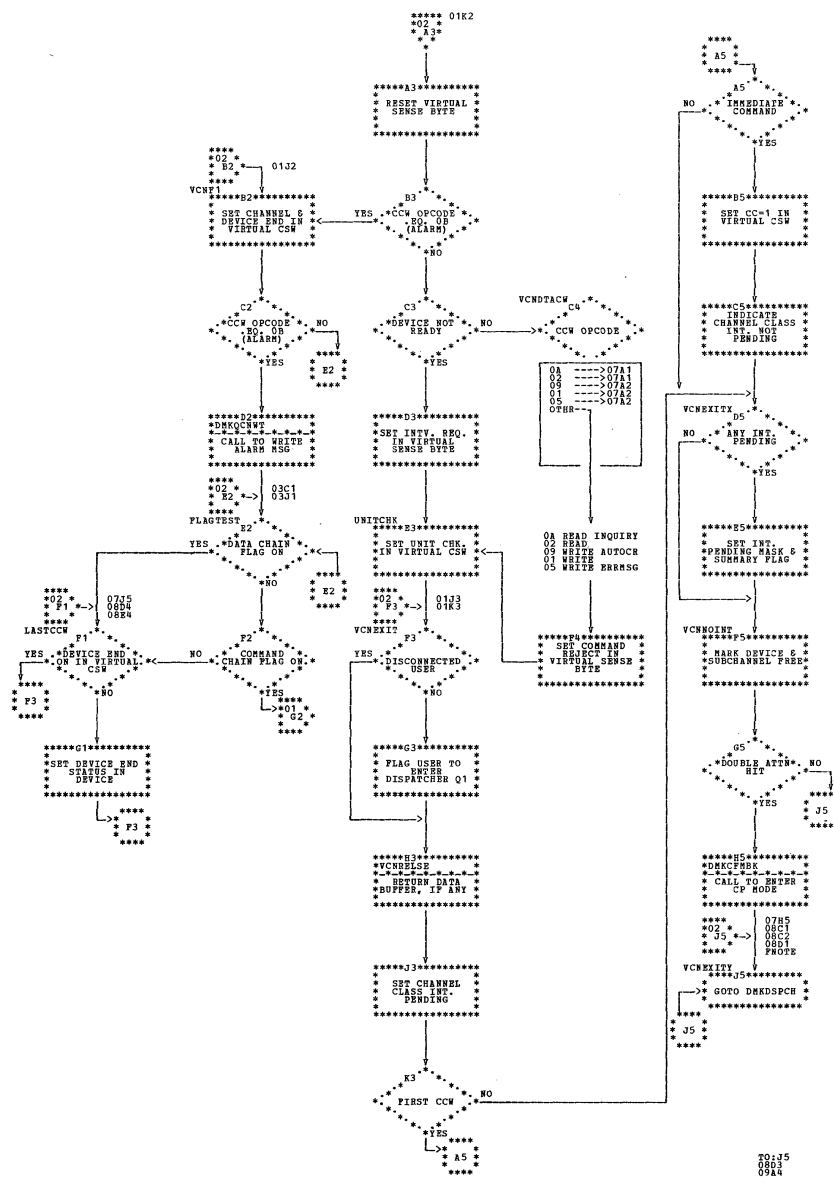


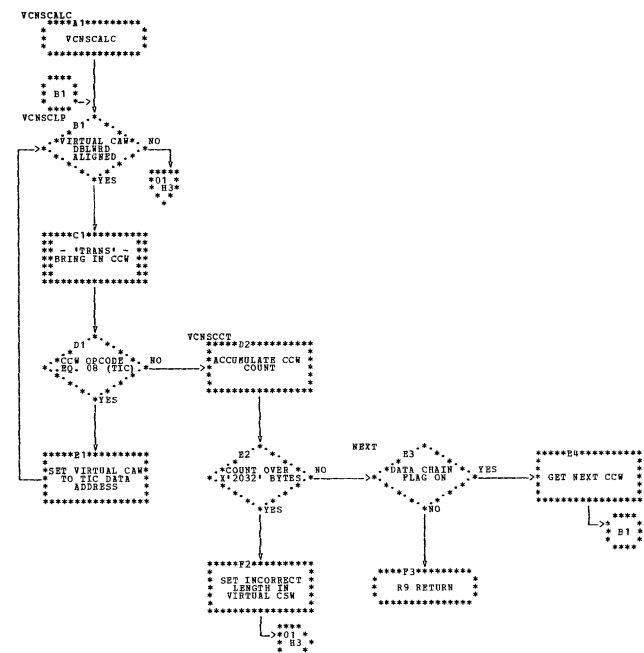
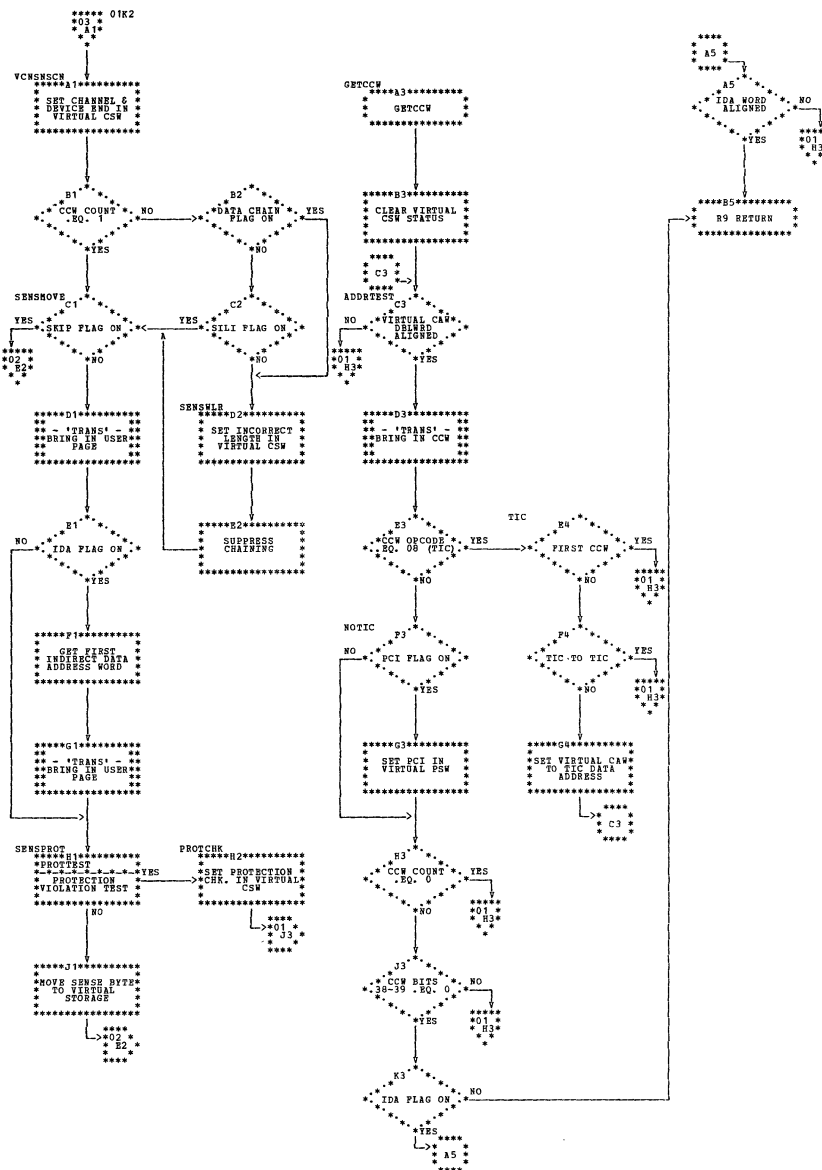


DMKVCN -- Console I/O Simulator (Parts 1 and 2 of 9)

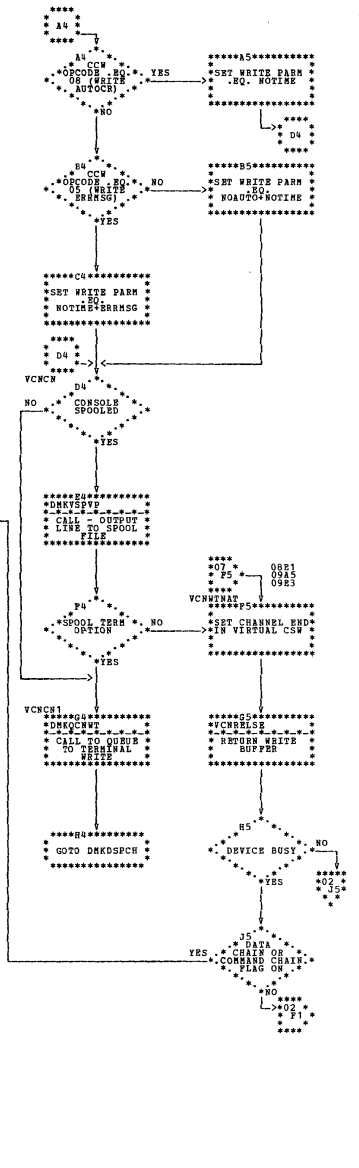
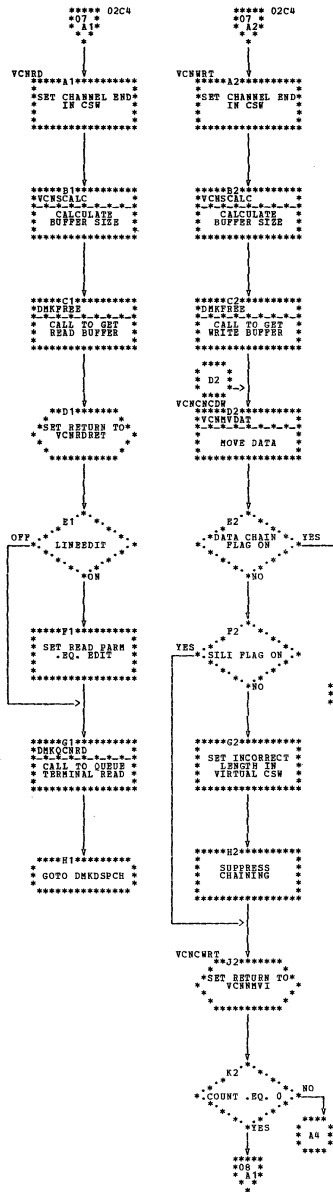


\* \* \* \* \*



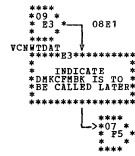
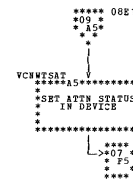
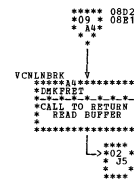
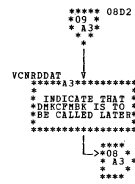
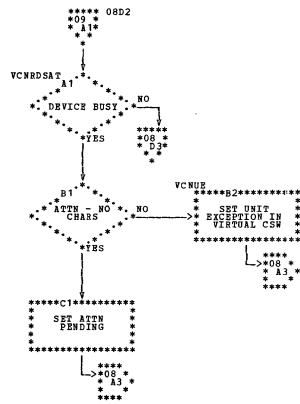






| DMKVCN -- Console I/O Simulator (Parts 7 and 8 of 9)

| DMKVCN -- Console I/O Simulator (Part 9 of 9)







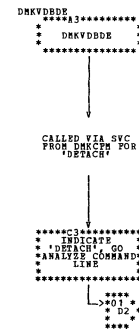
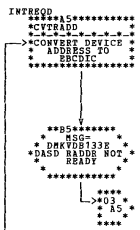
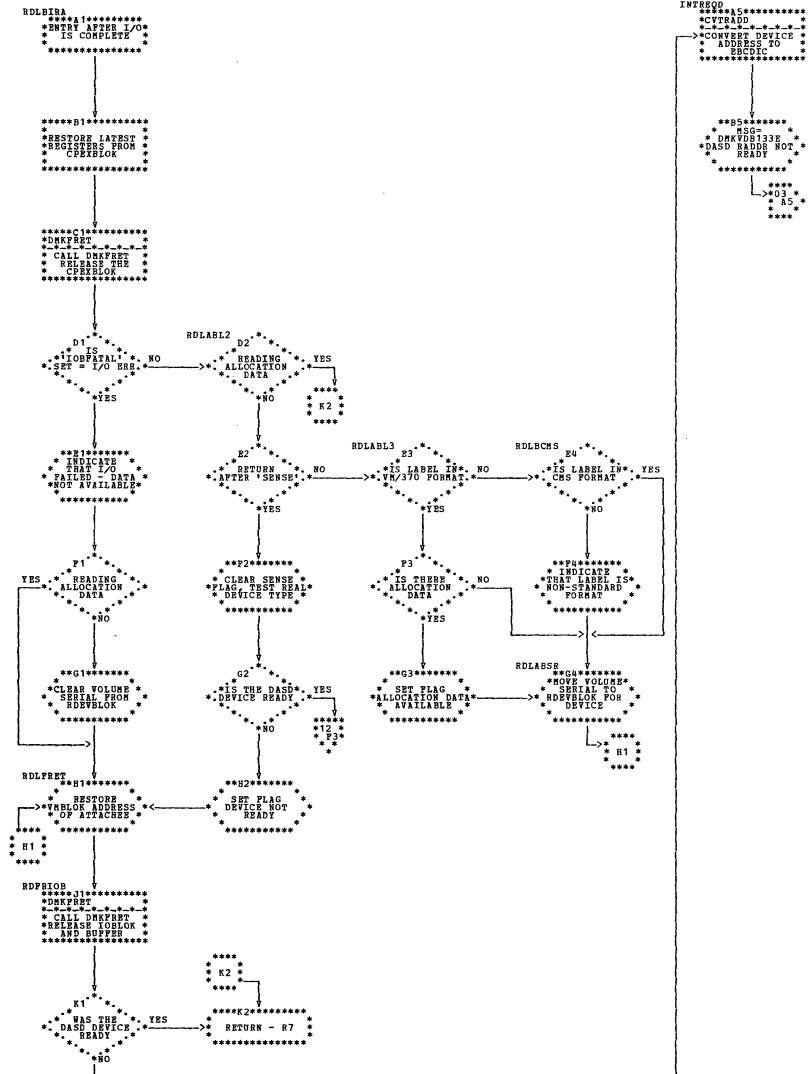








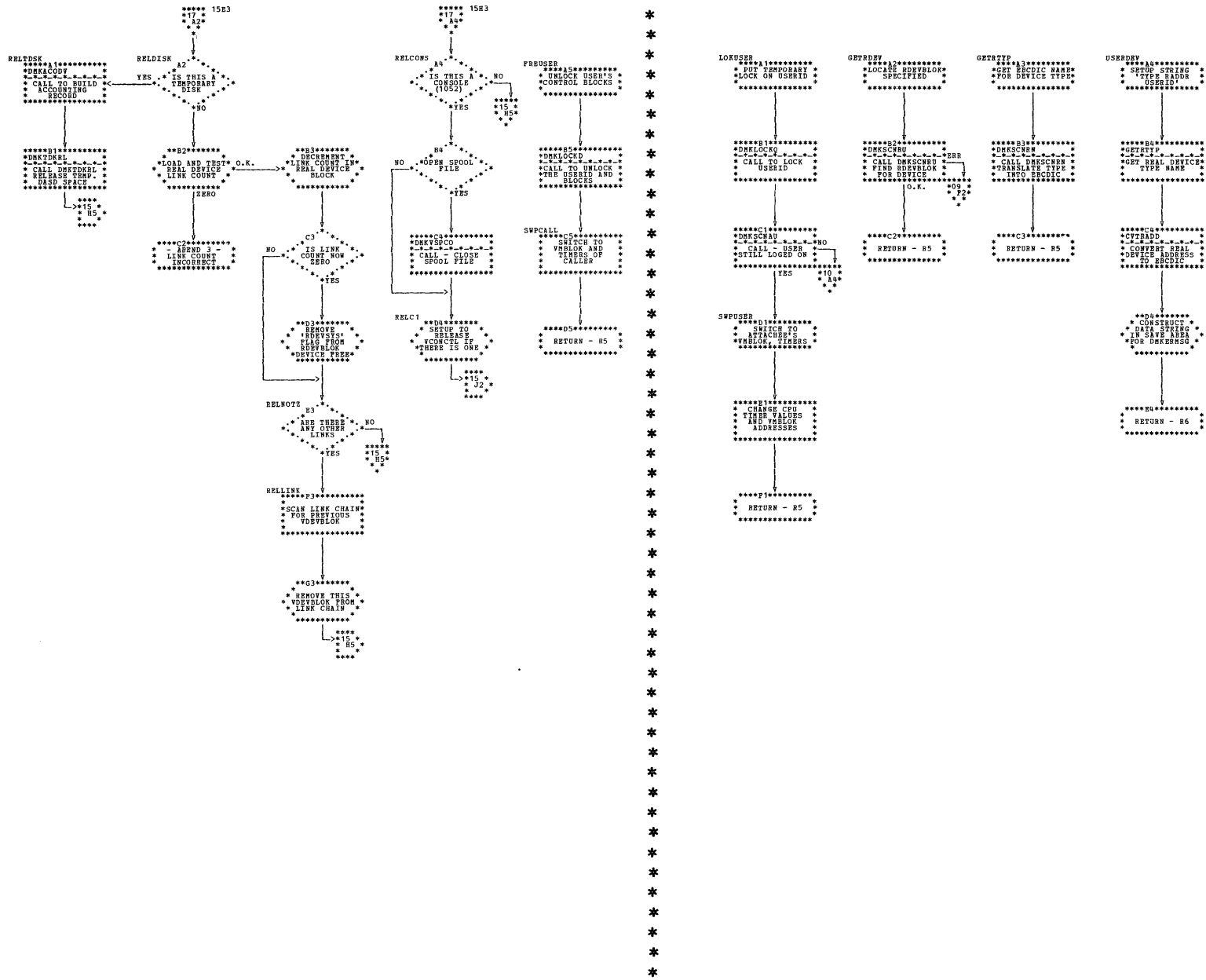




| DMKVDB -- Process ATTACH and DETACH Commands (Parts 13 and 14 of 18)

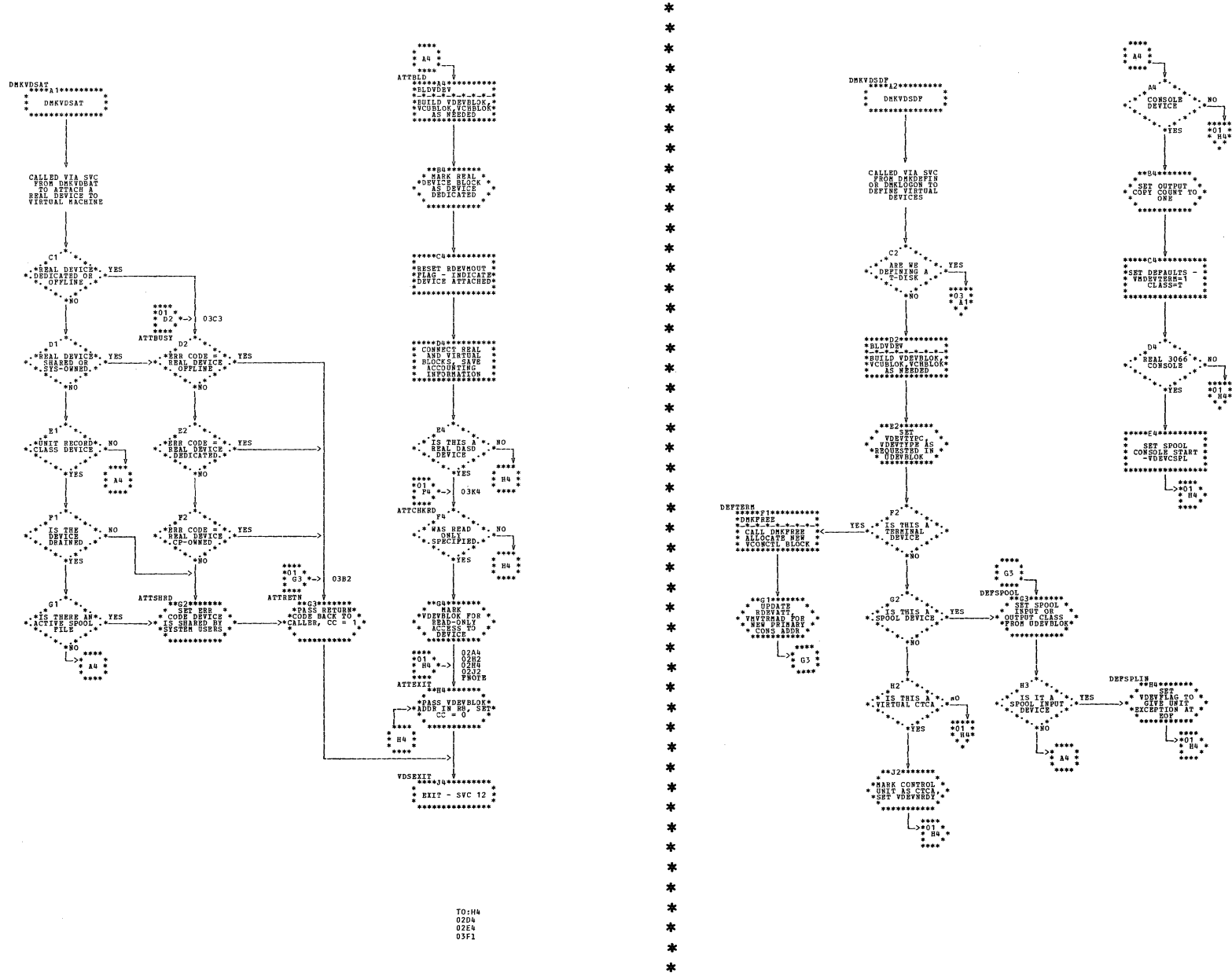






| DMKVDB -- Process ATTACH and DETACH Commands (Parts 17 and 18 of 18)

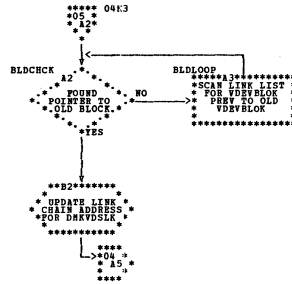
DMKVD S -- ATTACH, DEFINE, and LINK Virtual Device Subroutines (Parts 1 and 2 of 5)



TO:H4  
O2D4  
O2E4  
O3F1

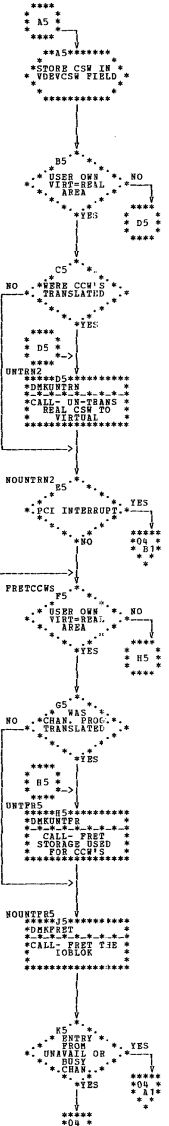
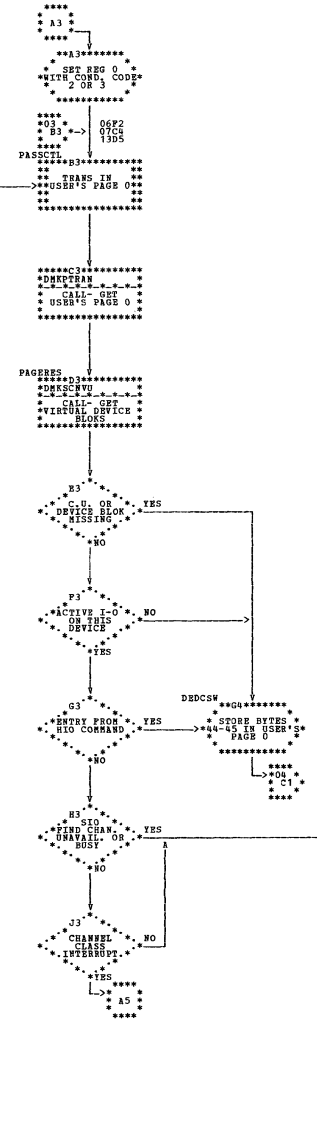
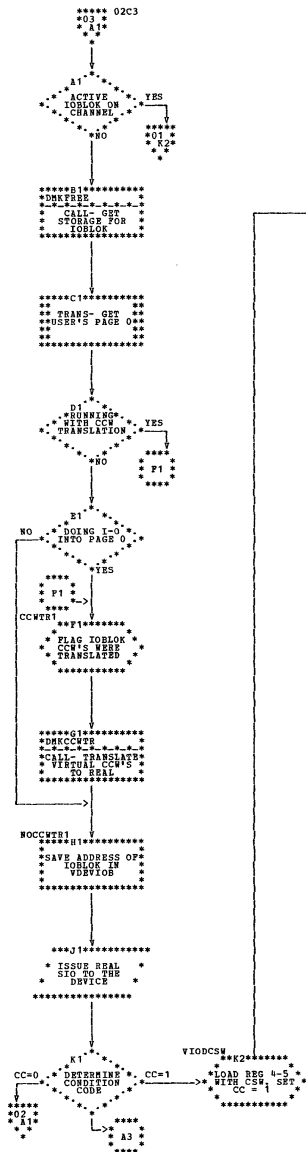


| DMKVDS -- ATTACH, DEFINE, and LINK Virtual Device Subroutines (Part 5 of 5)

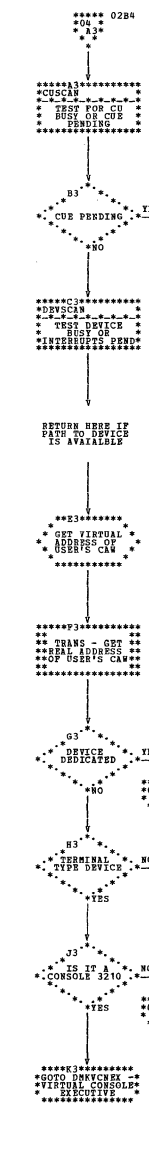
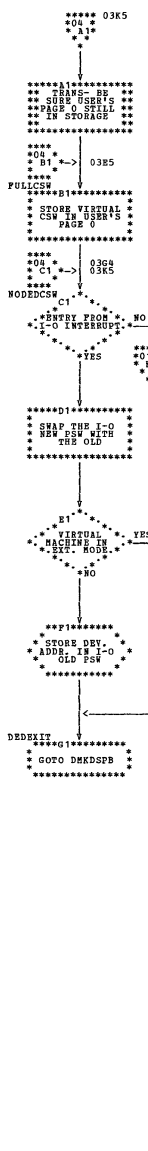




| DMKVIO -- Virtual I/O Manager (Parts 3 and 4 of 14)



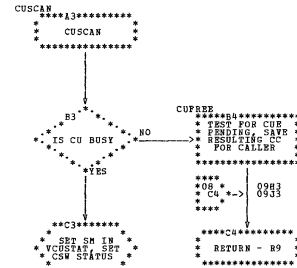
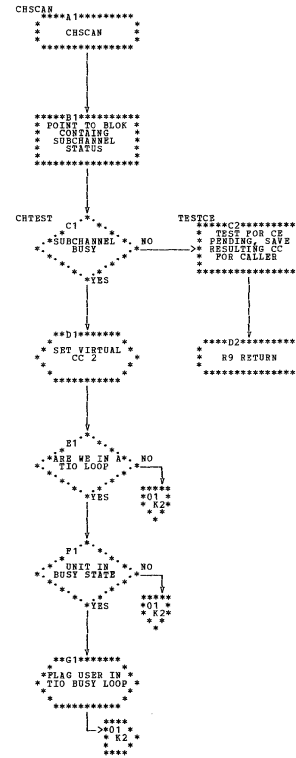
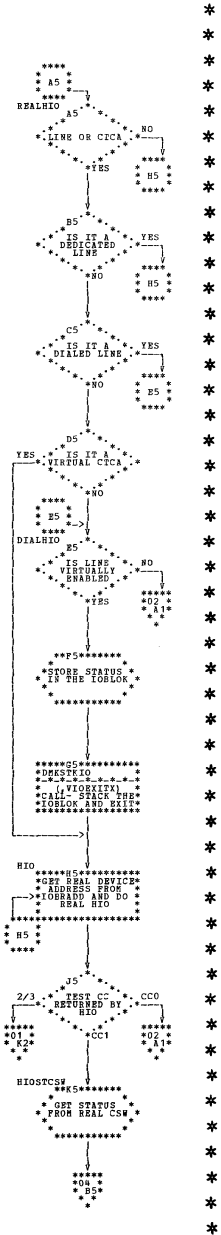
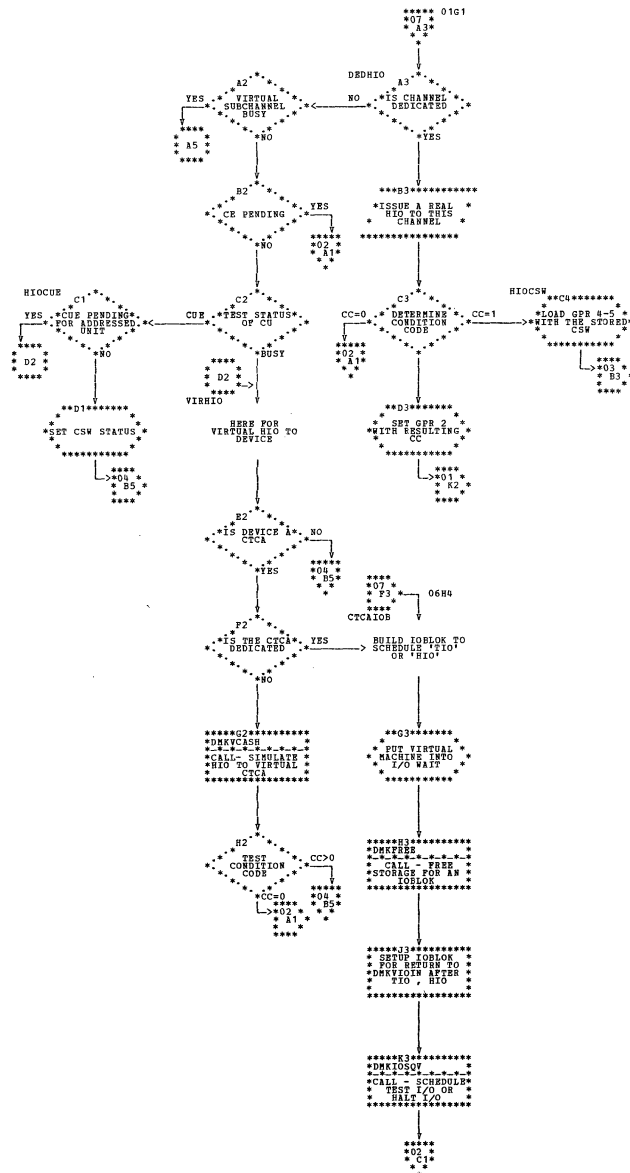
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| DMKVIC -- Virtual I/O Manager (Parts 7 and 8 of 14)



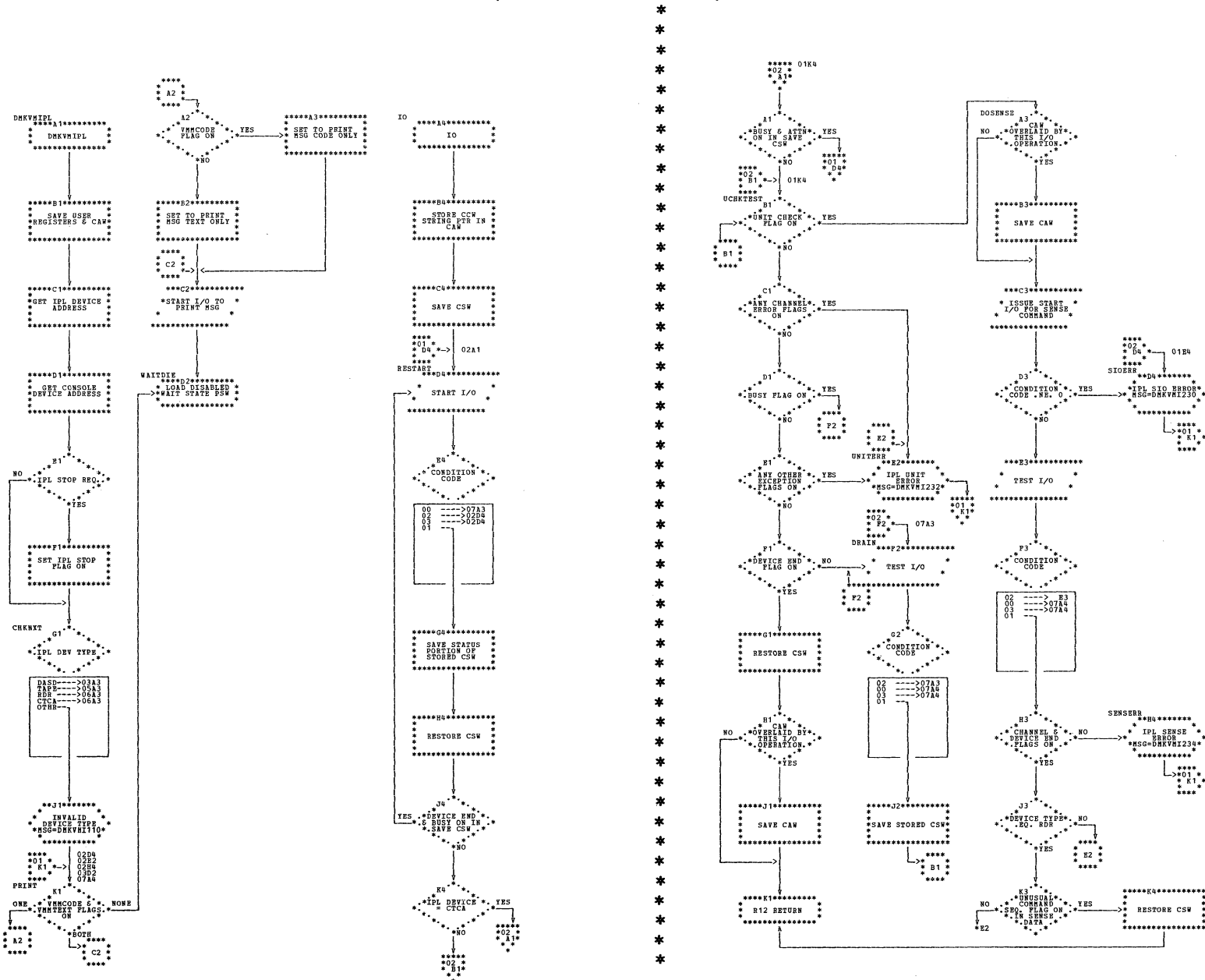


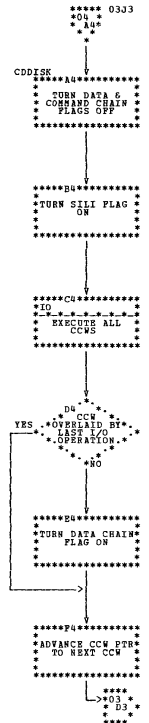
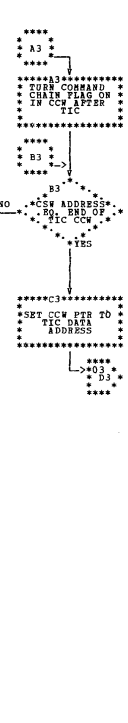
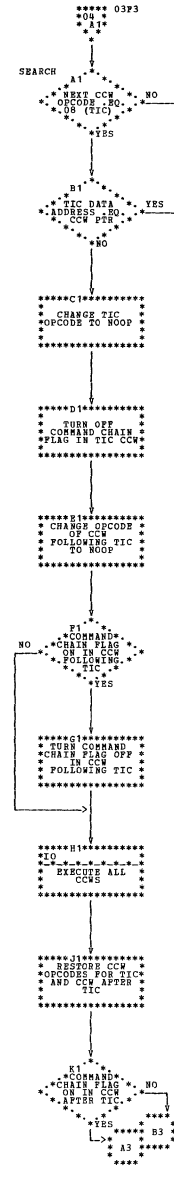
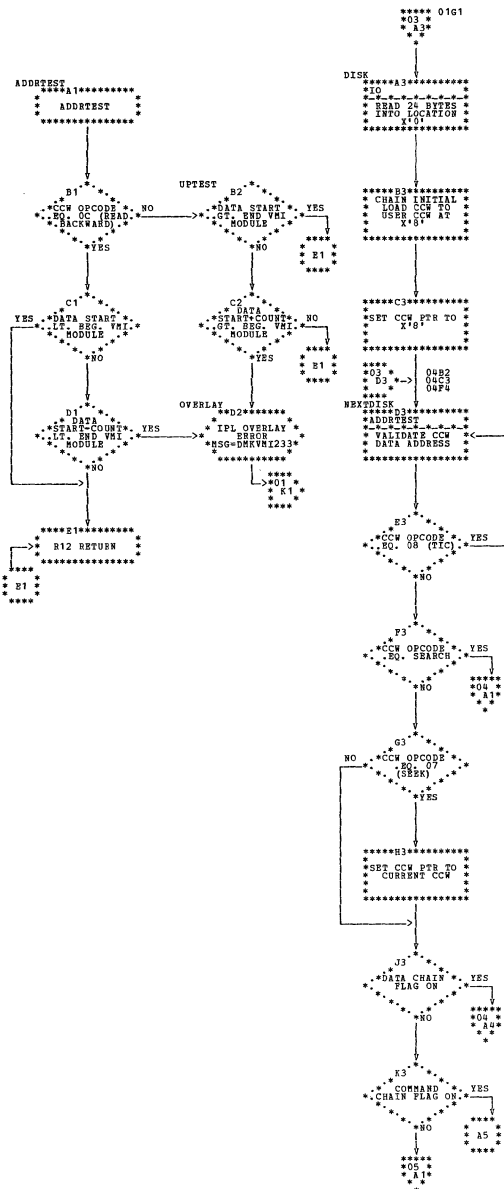






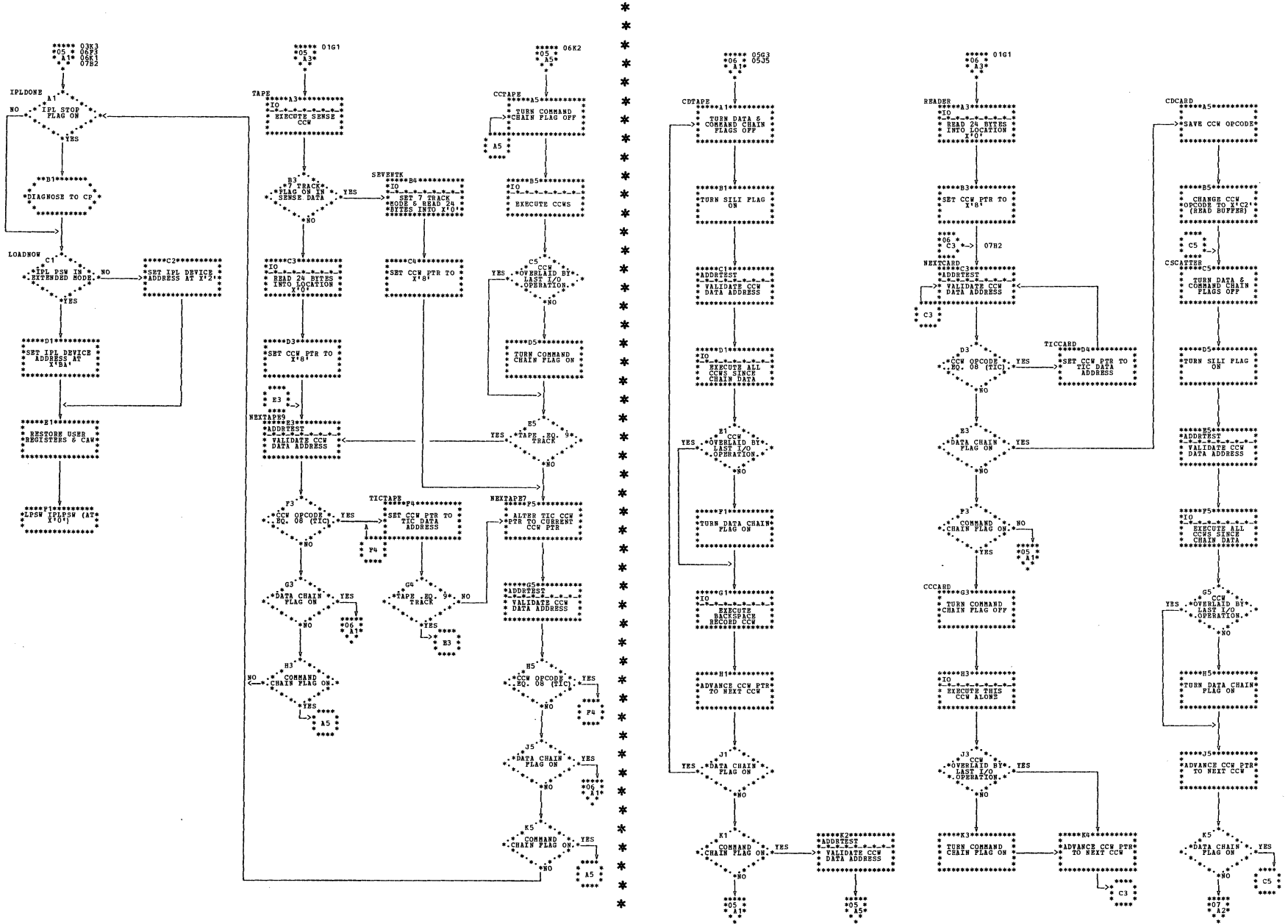
| DMKVMI -- Simulate IPL to Virtual Machine (Parts 1 and 2 of 7)

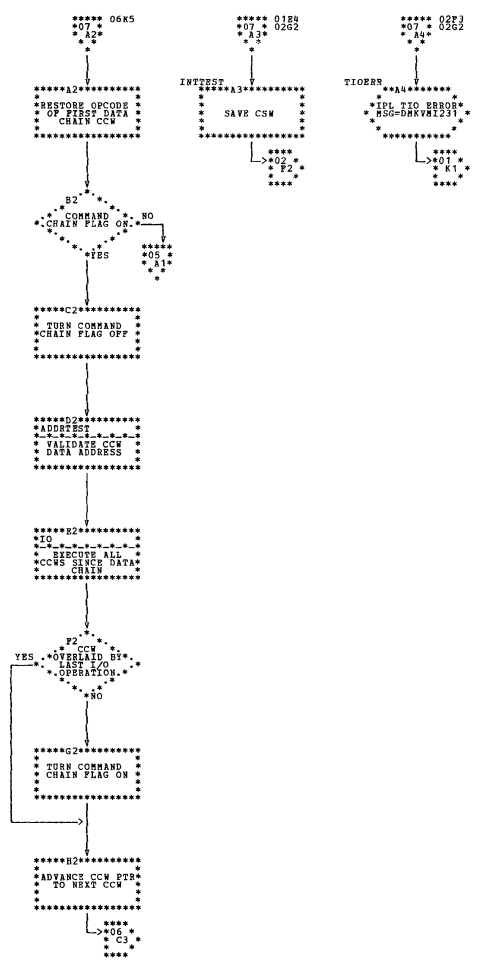




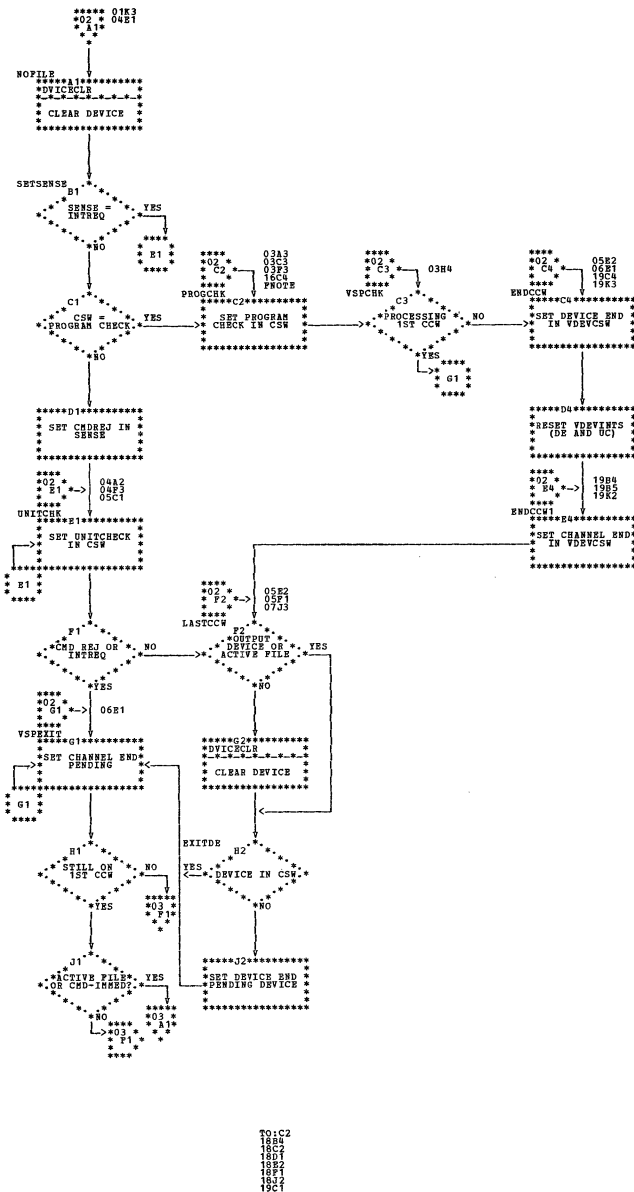
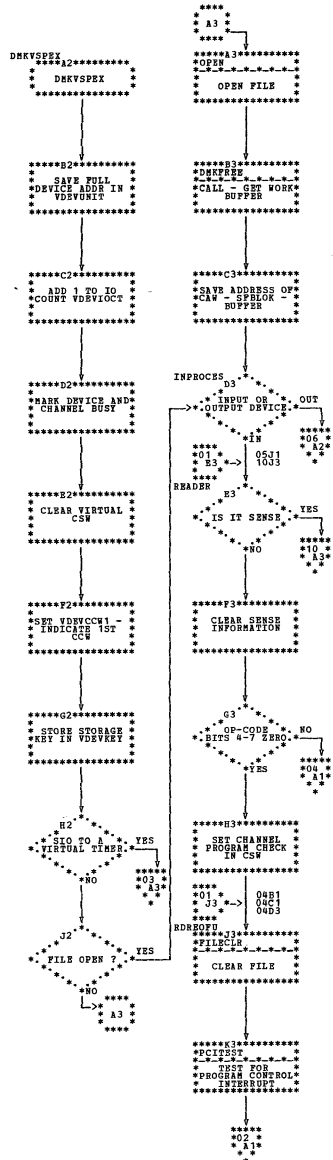
[ DMK VMI -- Simulate IPL to Virtual Machine (Parts 3 and 4 of 7)

| DMKVMI -- Simulate IPL to Virtual Machine (Parts 5 and 6 of 7)



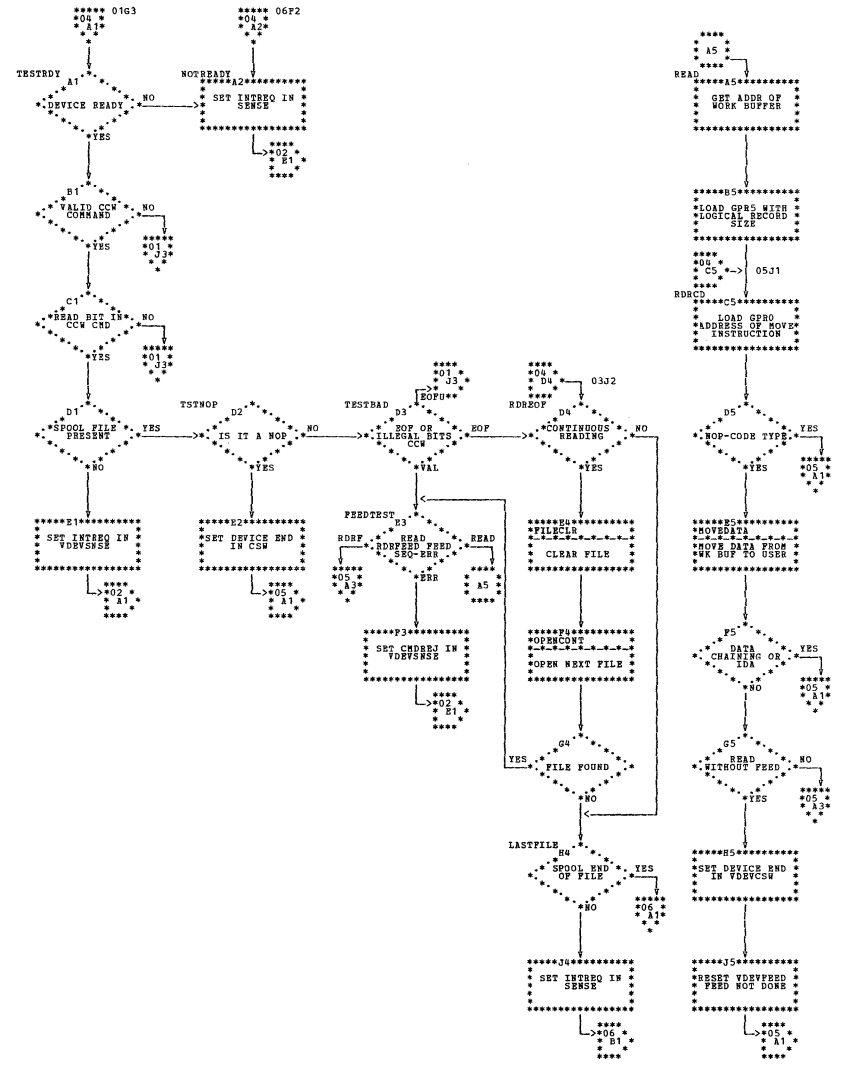
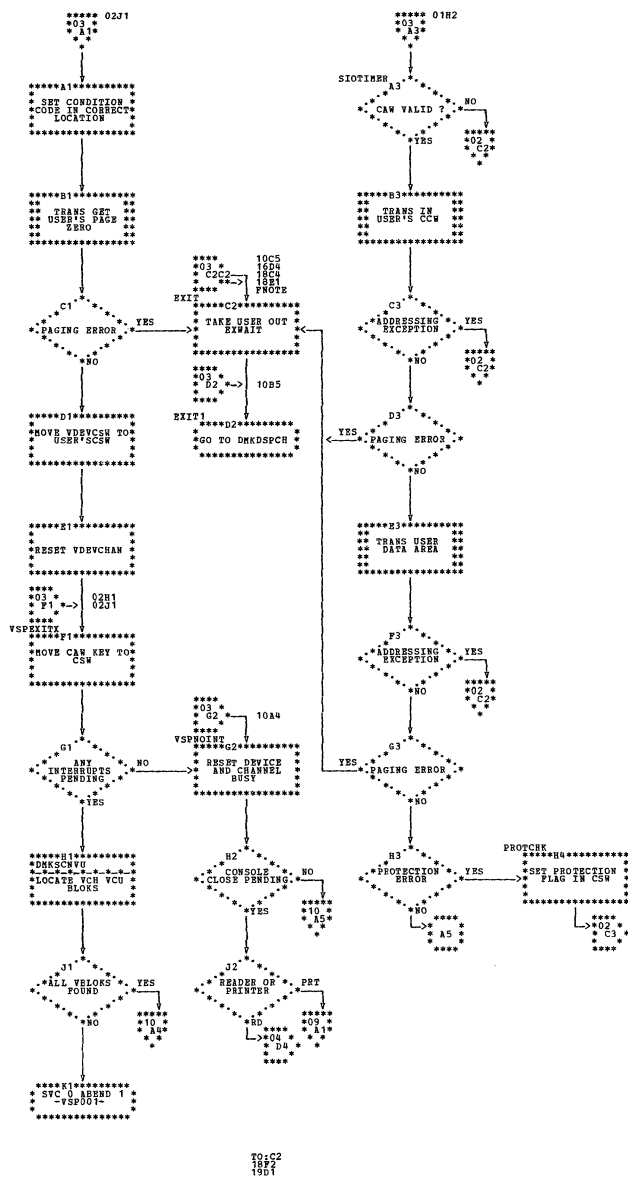


DMKVSP -- Virtual Spooling Manager (Parts 1 and 2 of 19)



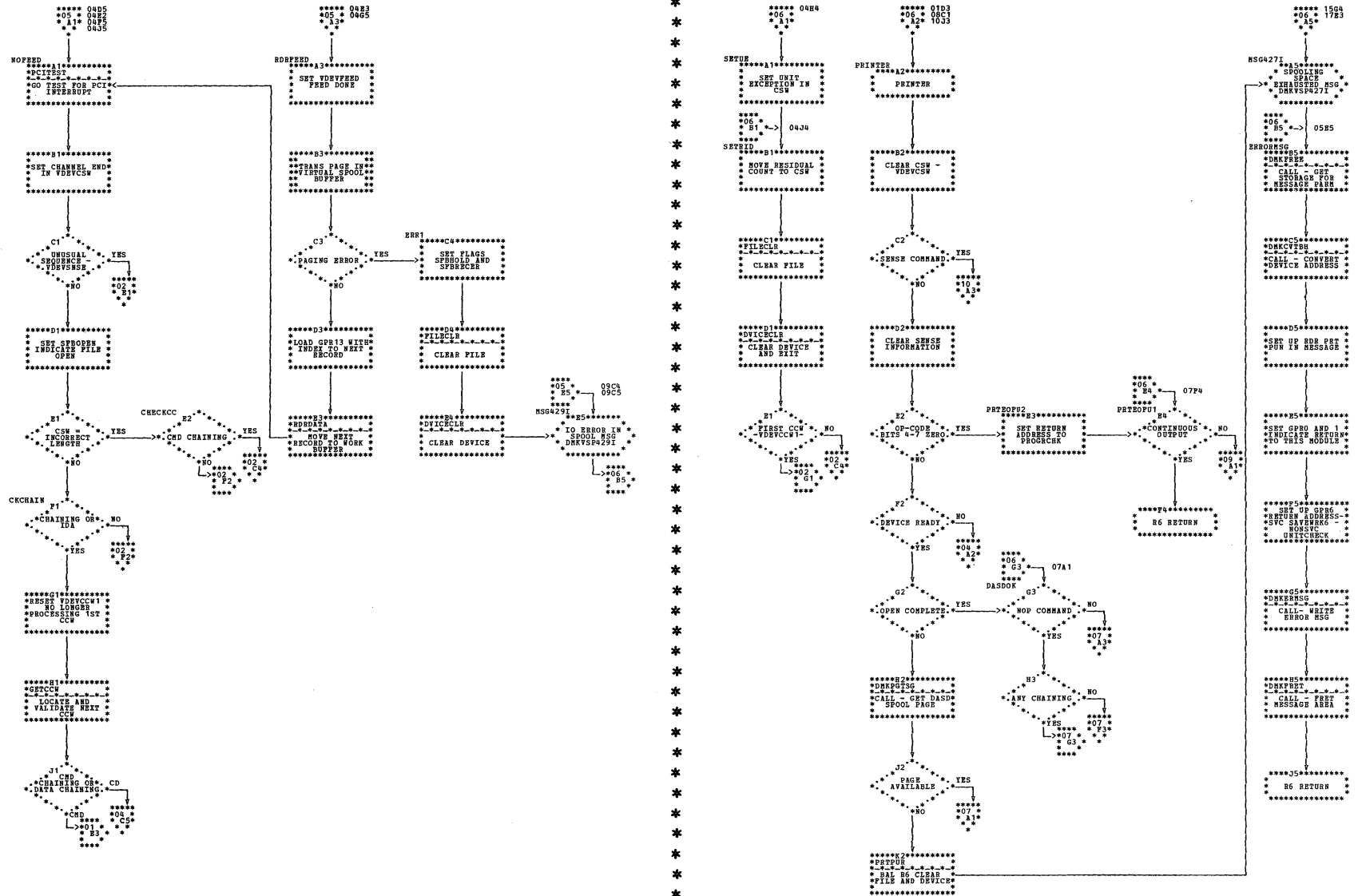
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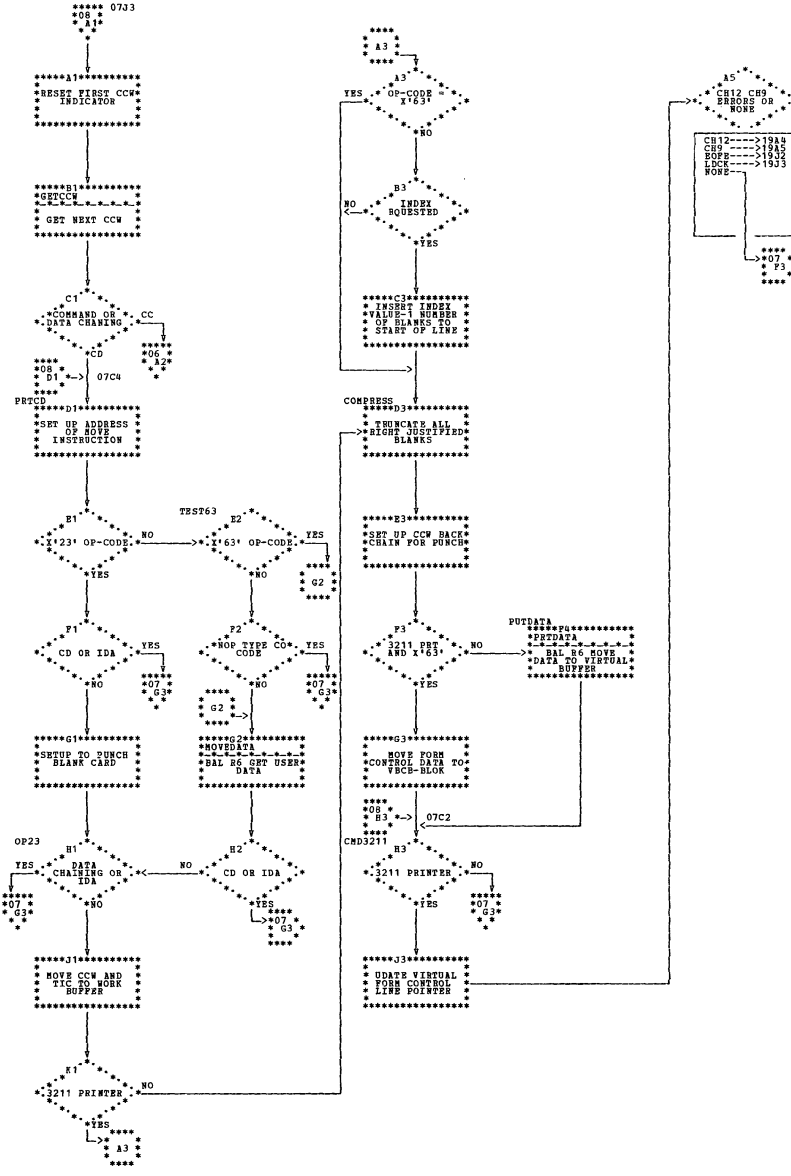
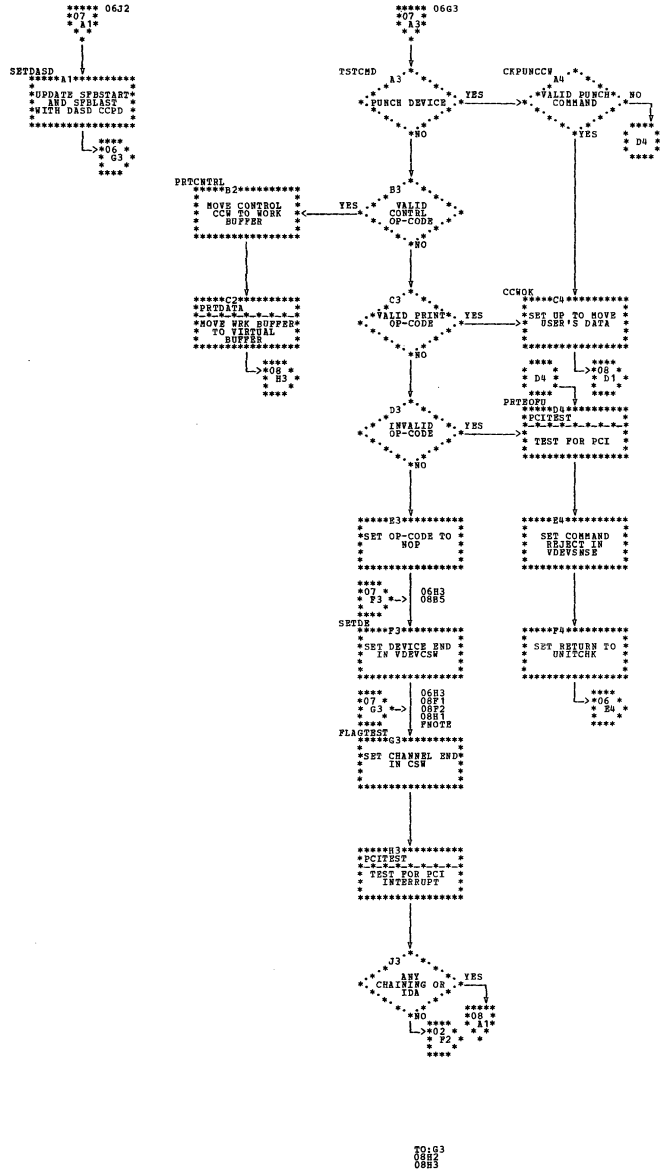




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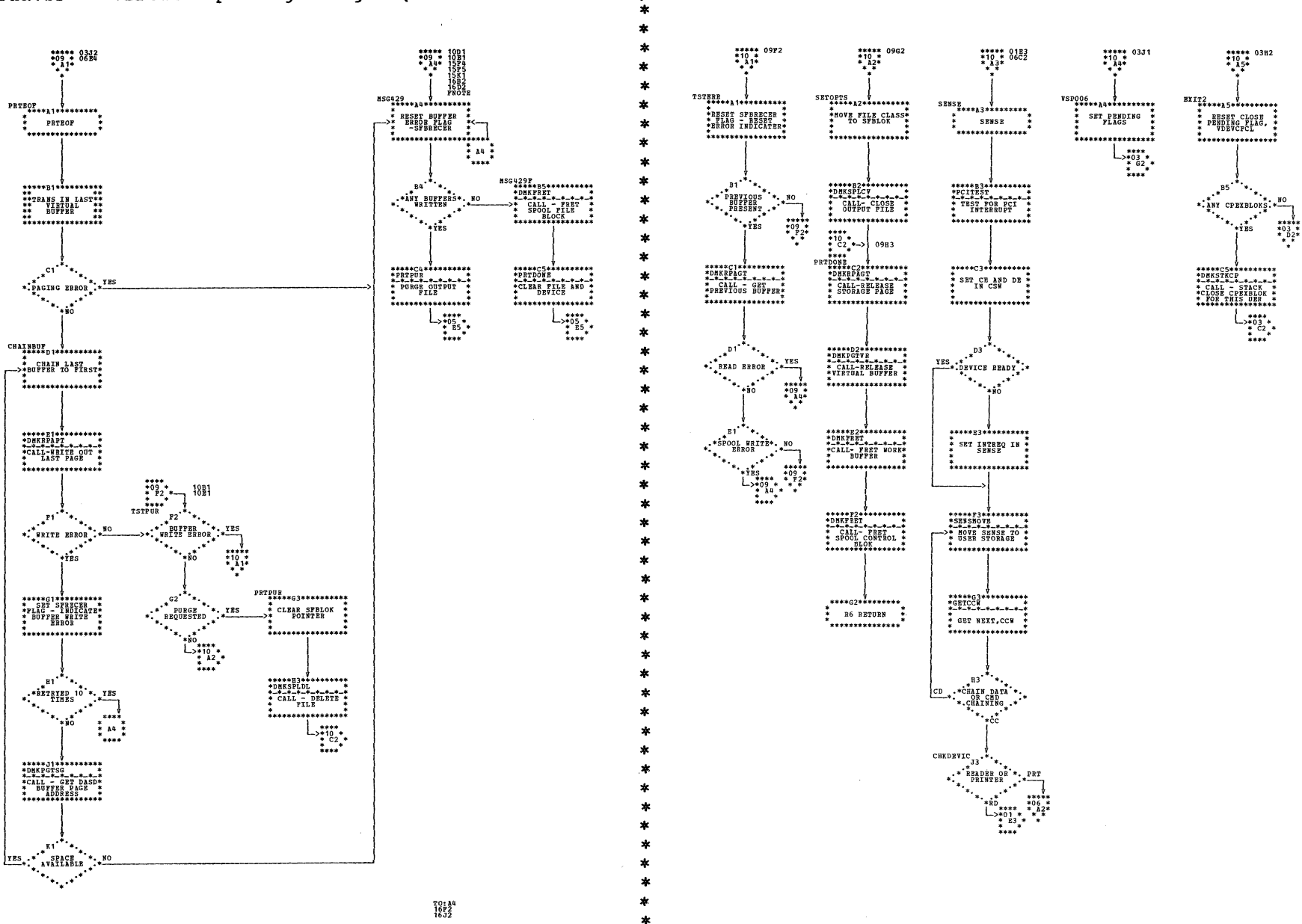
| DMKVSP -- Virtual Spooling Manager (Parts 5 and 6 of 19)





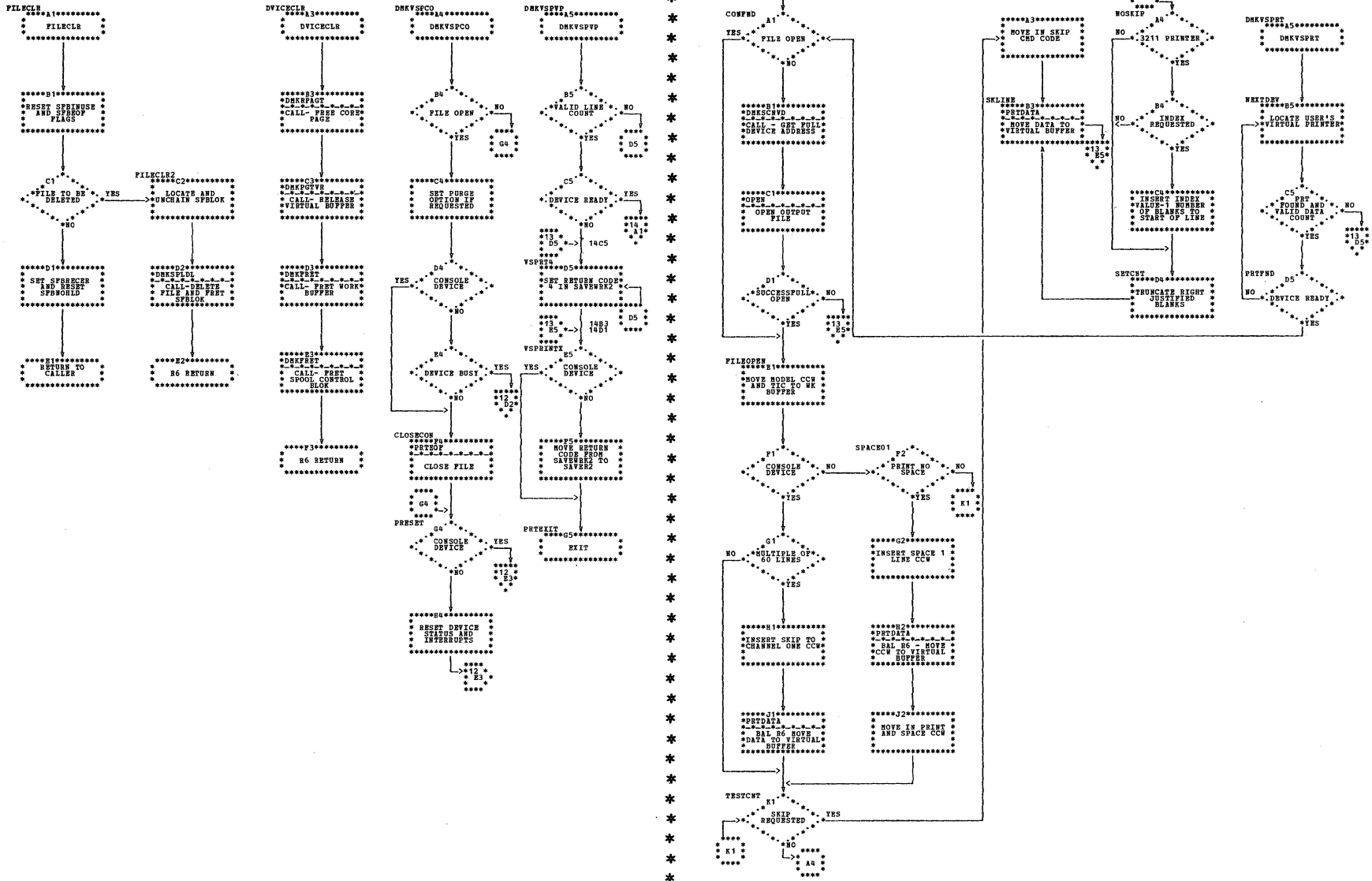
| DMKVSP -- Virtual Spooling Manager (Parts 7 and 8 of 19)

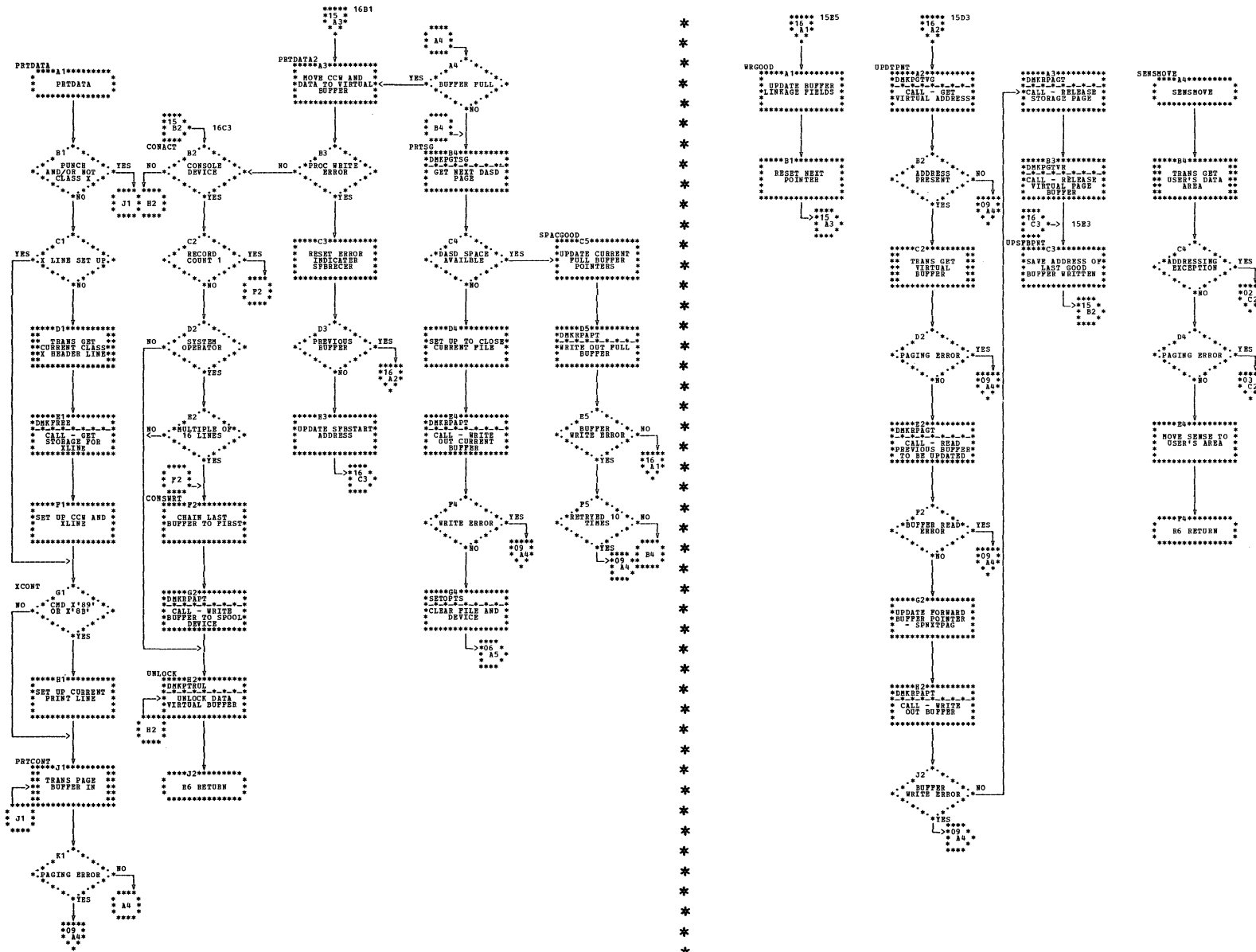
| DMKVSP -- Virtual Spooling Manager (Parts 9 and 10 of 19)





DMKVSP -- Virtual Spooling Manager (Parts 13 and 14 of 19)

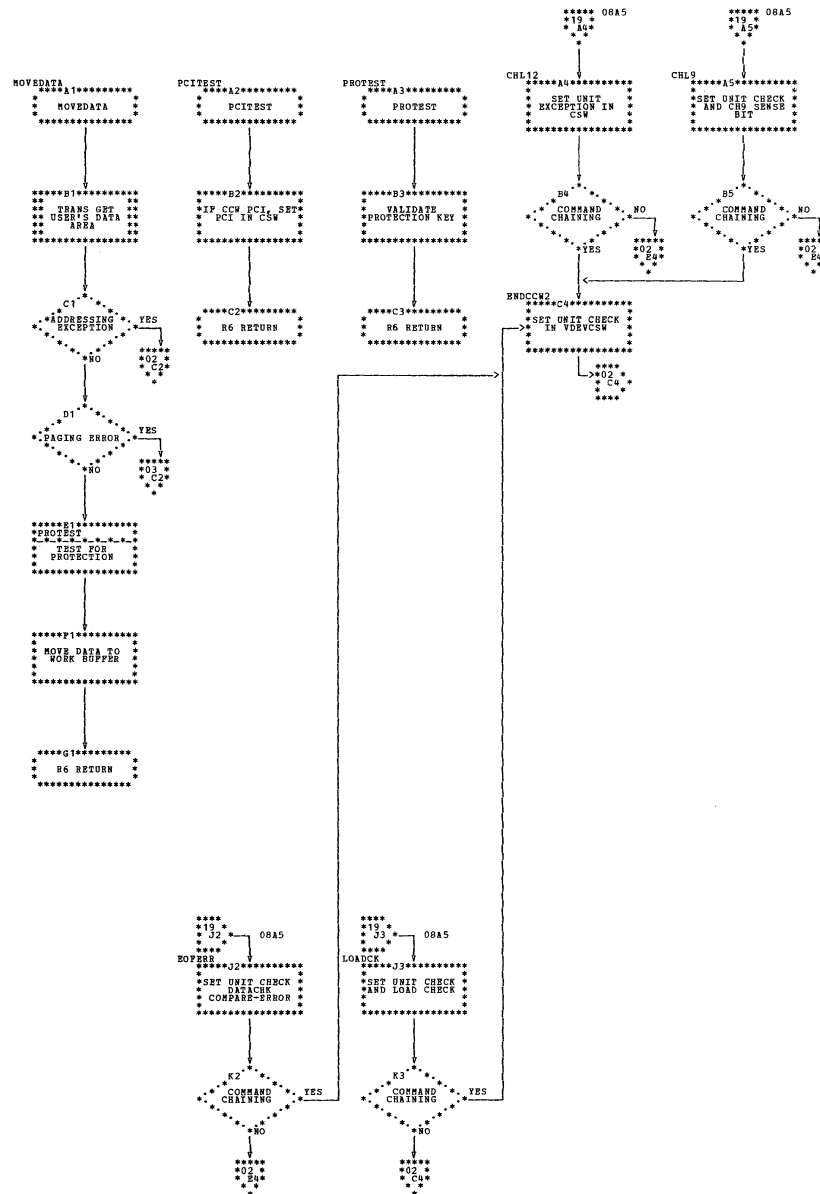


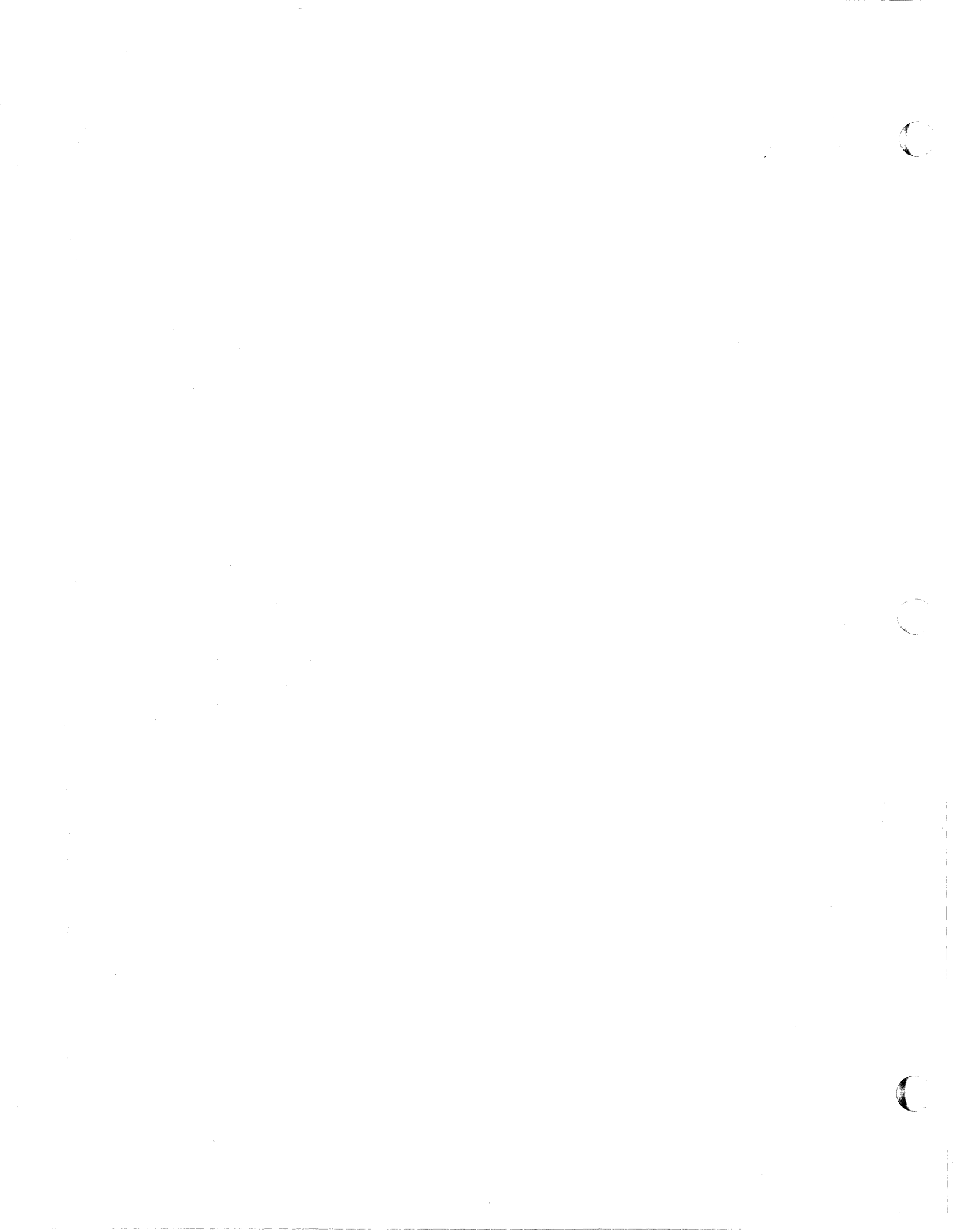


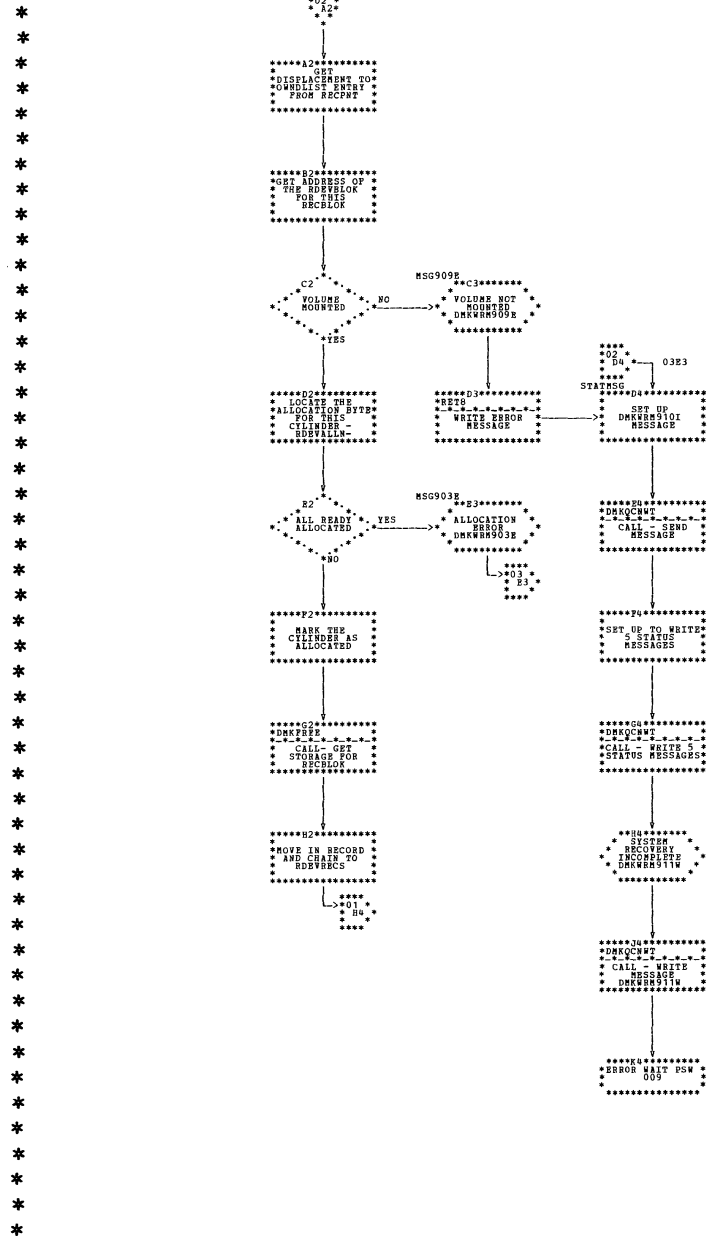
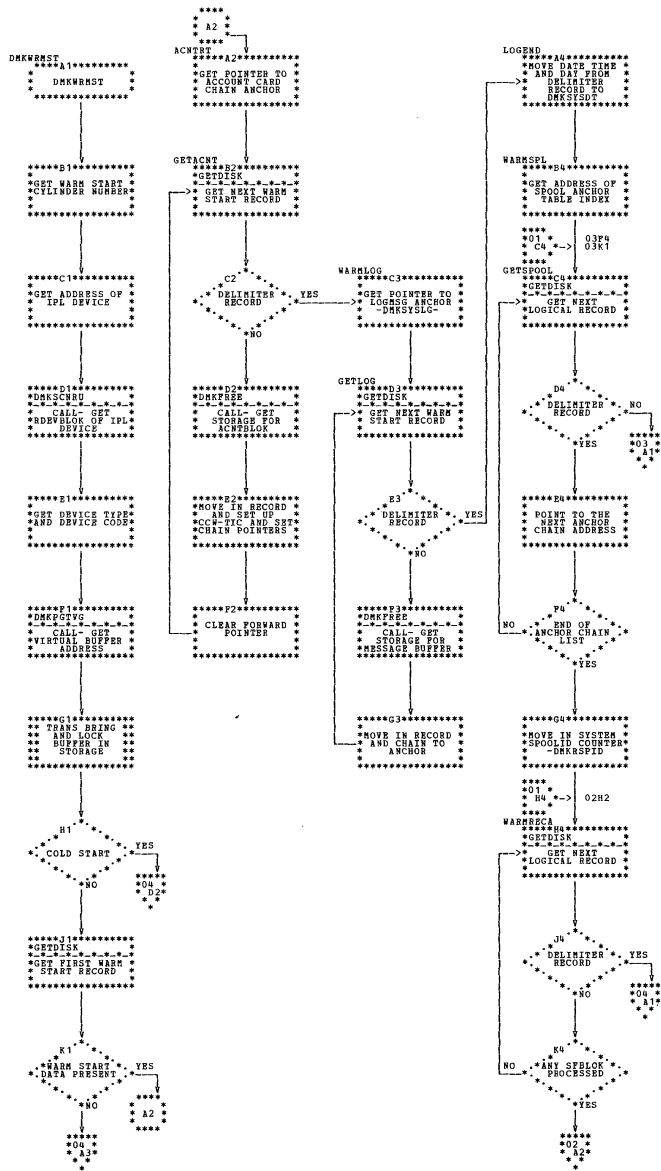
| DMKVSP -- Virtual Spooling Manager (Parts 15 and 16 of 19)







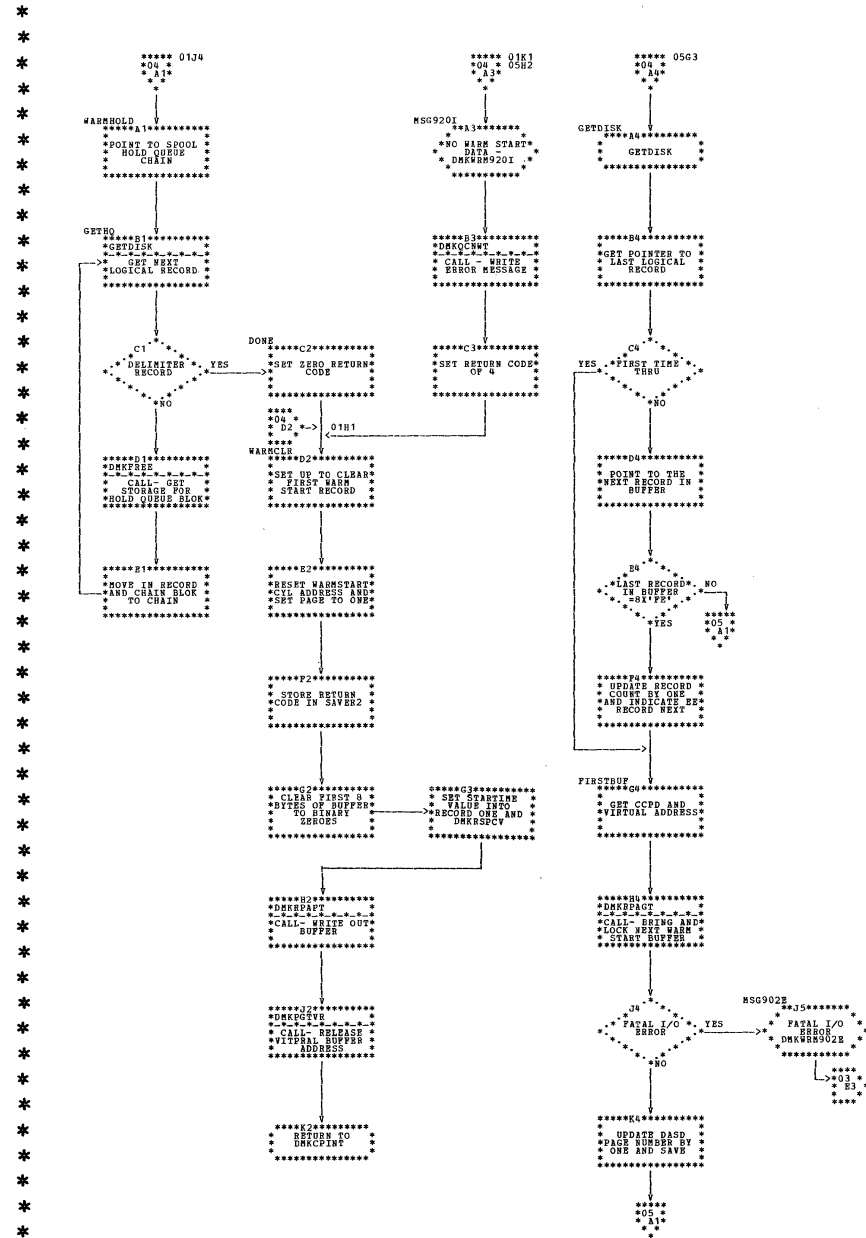
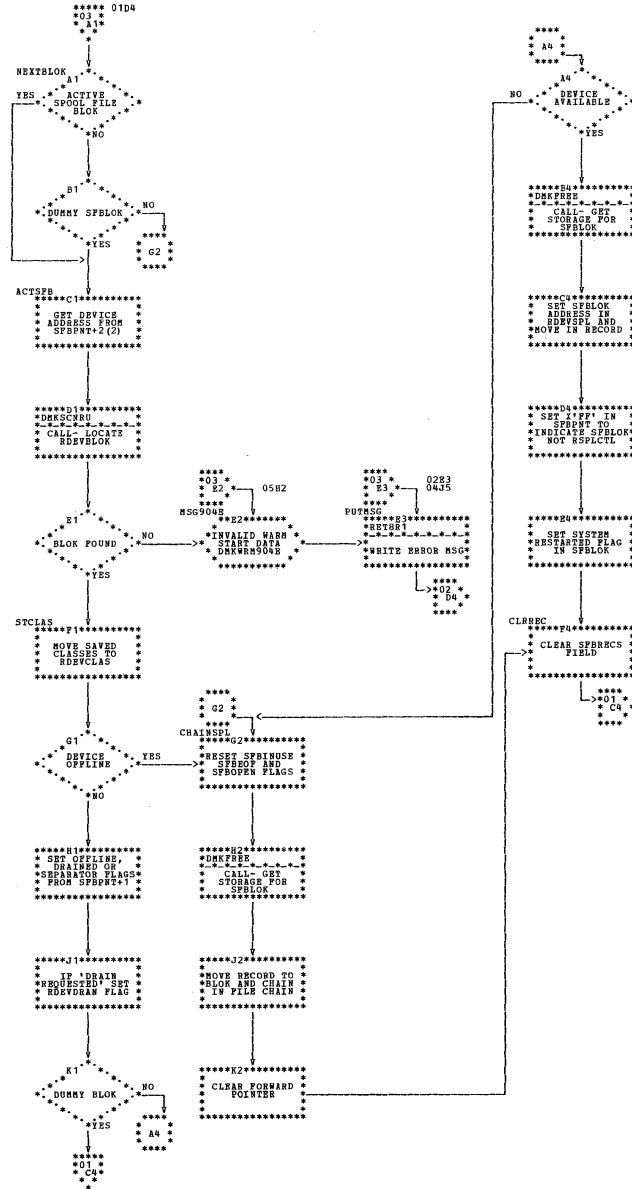




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DMKWRM -- Warm Start - (Parts 1 and 2 of 5)

DMKWRM -- Warm Start (Parts 3 and 4 of 5)







## DIRECTORIES

### MODULE/ENTRY POINT DIRECTORY

The directory contains much of the CP cross-reference information for the manual.

Name contains 1 of 2 possible items:

Modules - Six alphabetic characters that begin with DMK.

Entry Points - Seven or eight alphameric characters that begin with DMK. The fourth, fifth, and sixth characters of the entry point identify the module.

Comments Contains:

For modules, the function and a list of the data area that are used.

For entry points, the function.

MOD indicates:

Method of Operation Diagram that the entry point is in.

Chart indicates:

Page of the flowchart for a module that the entry point is.

Calls To indicates:

For modules the entry points in other modules that this module calls (Calls To).

Called By indicates:

For entry points, the modules that call it (Called By).

Name	Comments	MOD	Chart	Calls To	Called By
DMKACO	Accounting Routines			DMKCVTBH	
	<u>Data Areas Used</u>			DMKDSPCH	
	ACNTBLOK			DMKERMSG	
	IOBLOK			DMKFREE	
	RDEVBLOK			DMKFRET	
	SAVEAREA			DMKIOSQR	
	SYSLOCS			DMKPTRUL	
				DMKQCNWT	
				DMKRSPEX	

Name	Comments	MOD	Chart	Calls To	Called By
				DMKSTKCP	
				DMKSTKIO	
DMKACODV	Create account card for VDEVBLOK		1		DMKCPV DMKDIA DMKUSO DMKVDB
DMKACOFF	Create account card for VMBLOK		1		DMKCPV
DMKACON	User entry point		2		DMKLOG DMKUSO
DMKACOPU	Punch queued account cards		3		DMKRSP
DMKACOQU	Queue request to punch cards		3		
DMKACOTM	Process use time message		2		DMKCPV DMKCQG DMKUSO
DMKBLD	Process Blocks <u>Data Areas Used</u> PGTABLE RDEVBLOK SEGTABLE SWPTABLE VMBLOK			DMKCVTBH DMKERMSG DMKFREE DMKFRET DMKSCNRD	
DMKBLDRL	Release real segment, page, and swap tables.		2		DMKDEF DMKUSO
DMKBLDRT	Build real segment, page, and swap tables.		1		DMKCNS DMKCPI DMKDEF DMKLOG
DMKBLDVM	Build a new VMBLOK		3		DMKCNS DMKDIA
DMKCCH	Channel Check Handler <u>Data Area Used:</u> CCHREC IOBLOK IOERBLOK PSA RCHBLOK RCUBLOK RDEVBLOK SAVEAREA VMBLOK			DMKCVTBH DMKFREE DMKIOECC DMKQCNWT DMKSCNRU	
DMKCCHIS	Channel check and CSW stored after SIO.	8B2	1		DMKIOS
DMKCCHNT	Channel check from I/O interrupt		3		DMKIOS
DMKCCHRT	Print error message.		3		DMKIOE



Name	Comments	MOD	Chart	Calls To	Called By
DMKCCW	Translate CCWs <u>Data Areas Used:</u> IOBLOK RDEVBLOK SAVEAREA VDEVBLOK VMBLOK				DMKDIASM DMKFREE DMKFRET DMKISMTR DMKPTRAN DMKPTRLK DMKUNTRS
DMKCCWSB	Get SEEK arguments.		30		DMKTRC
DMKCCWTR	Translate user's CCWs.		1		DMKGEN DMKVIO
DMKCDB	Process DISPLAY, DCP, DUMP, and DMCP commands. <u>Data Areas Used</u> BUFFER ECBLOK PGTABLE SAVEAREA VMBLOK				DMKCVTBD DMKCVTBH DMKCVTDB DMKCVTFP DMKCVTHB DMKDSPCH DMKERMSG DMKFREE DMKFRET DMKPTRAN DMKQCNWT DMKSCNFD DMKVATAB DMKVATBC DMKVATMD DMKVSPT
DMKDBDC	Display real storage (DCP Command).		4		DMKCFM
DMKDBDI	Display virtual storage (DISPLAY command).		1		DMKCFM
DMKDBDM	Dump real storage in spooled printer (DMCP command).		4		DMKCFM
DMKDBDU	Dump virtual storage on spooled printer (DUMP command).		3		DMKCFM
DMKCDS	Process STORE and STCP commands. <u>Data Areas Used:</u> ECBLOK SAVEAREA TRQBLOK VMBLOK				DMKCVTBH DMKCVTDB DMKCVTHB DMKERMSG DMKPTRAN DMKQCNWT DMKSCNFD DMKVATAB DMKVATBC

Name	Comments	MOD	Chart	Calls To	Called By
DMKCDSCP	Store into real storage (STCP command).		1	DMKVATMD	DMKCFM
DMKCDSTO	Store into virtual storage (STORE command).		1		DMKCFM
DMKCFD	Process LOCATE and ADSTOP commands. <u>Data Areas Used:</u> ECBLOK RCHBLOK RDEVBLOK SAVEAREA VCHBLOK VCUBLOK VDEVBLOK VMBLOK			DMKCVTBH DMKCVTHB DMKERMSG DMKFREE DMKFRET DMKQCNWT DMKSCNAU DMKSCNFD DMKSCNRU DMKSCNVU	
DMKCFDAD	Stop virtual machine at specified address (ADSTOP command).		3		DMKCFM
DMKCFDLO	Display address of real device blocks, or VMBLOK and/or virtual device blocks (LOCATE command).		1		DMKCFM
DMKCFG	Command Processor <u>Data Areas Used:</u> ECBLOK PSA RDEVBLOK SAVEAREA SAVTABLE SYSTABLE VDEVBLOK VMBLOK			DMKCVTBH DMKERMSG DMKFREE DMKFRET DMKPTRAN DMKPTRUL DMKQCNRD DMKQCNWT DMKRPAPT DMKSCNFD DMKSCNVS DMKSCNVU	
DMKCFGSV	SAVESYS command processor.		1		DMKCFM
DMKCFM	Process SLEEP, BEGIN and QUERY command and direct all other CP commands to current module. <u>Data Areas Used:</u> BUFFER SAVEAREA VMBLOK			DMKCDBDC DMKCDBDI DMKCDBDM DMKCDBDU DMKCDSCP DMKCDSTO DMKCFDAD DMKCFDLO DMKCFMBE	

Name	Comments	MOD	Chart	Calls To	Called By
DMKCFM (cont)				DMKCFMQU	
				DMKCFMSL	
				DMKCFPEX	
				DMKCFPIP	
				DMKCFPNR	
				DMKCFPRS	
				DMKCFPRW	
				DMKCFPRY	
				DMKCFPSR	
				DMKCFSET	
				DMKCFTRM	
				DMKCPVAC	
				DMKCPVDS	
				DMKCPVEN	
				DMKCPVH	
				DMKCPVLK	
				DMKCPVRY	
				DMKCPVSH	
				DMKCPVSV	
				DMKCPVUL	
				DMKCQGEN	
				DMKCQPRV	
				DMKCSOBS	
				DMKCSODR	
				DMKCSOFL	
				DMKCSOLD	
				DMKCSORP	
				DMKCSOSP	
				DMKCSOST	
				DMKCSOVL	
				DMKCSPCL	
				DMKCSPPR	
				DMKCSPHL	
				DMKCSPSP	
				DMKCSUCH	
				DMKCSUOR	
				DMKCSUPU	
				DMKCSUTR	
				DMKCVTHB	
				DMKDEFIN	
				DMKDIAL	
				DMKDSPCH	
				DMKERMSG	
				DMKFREE	
				DMKFRET	
				DMKLNKIN	

Name	Comments	MOD	Chart	Calls To	Called By
DMKCFM (cont)					DMKLOGON DMKMCCCL DMKMSGEC DMKMSGMS DMKMSGWN DMKQCNRD DMKQCNWT DMKSCNFD DMKSCNVU DMKSTKCP DMKTRACE DMKUSODS DMKUSOFL DMKUSOLG DMKVDBAT DMKVDBDE
DMKCFMAT	Simulate attention interrupt to virtual machine.		3		DMKCNS
DMKCFMBE	BEGIN command processor		4		DMKCFM
DMKCFMBK	Attention interrupt twice from a terminal.	7B3	1		DMKCNS DMKDSP DMKHVC DMKIOE DMKMCH DMKPRG DMKPSA DMKPTR DMKTRC DMKVCN
DMKCFMEN	DIAGNOSE code 8.		2		DMKCP DMKHVC DMKLNK
DMKCFMQU	QUERY command processor (initial)		5		DMKCFM
DMKCFMSL	SLEEP command processor.		4		DMKCFM
DMKCFP	Simulate the operators console for the virtual machine. <u>Data Areas Used</u> IOBLOK RDEVBLOK SAVEAREA VCHBLOK VCUBLOK VDEVBLOK				DMKCVTBH DMKCVTDB DMKCVTHB DMKDSPCH DMKDIADR DMKERMSG DMKFREE DMKFRET DMKPGSPO

Name	Comments	MOD	Chart	Calls To	Called By
	VMBLOK			DMKPTRAN DMKPTRUL DMKRPAGT DMKQCNWT DMKSCHRT DMKSCNFD DMKSCNVS DMKSCNVU DMKSTKCP DMKSTKIO DMKUNTPR DMKVATBC DMKVCARD DMKVDBRL DMKVSPCO DMKVSPCR	
DMKCFPII	IPL from LOGON		5		DMKLOG
DMKCFPIP	IPL command processor.	7B3.1	9		DMKCFM
DMKCFPRD	Reset a virtual device.		2		DMKDEF DMKVDB
DMKCFPRR	Process system resets from other routines		1		DMKDEF DMKMCH DMKUSO
DMKCFSS	Process SET command <u>Data Areas Used:</u> CORTABLE IRMBLOK SAVEAREA RDEVBLOK VMBLOK			DMKCVTBH DMKCVTDB DMKCVTDT DMKCVTHB DMKERMSG DMKFREE DMKFRET DMKMCHMS DMKQCNRD DMKQCNWT DMKSCHRT DMKSCNAU DMKSCNFD DMKSCNRU DMKSTKIO	
DMKCFSET	SET command processor		1		DMKCFM
DMKCFST	Process user's terminal options. <u>Data Areas Used:</u>			DMKCVTDB DMKERMSG	

Name	Comments	MOD	Chart	Calls To	Called By
	RDEVBLK			DMKQCNWT	
	SAVEAREA			DMKSCNFD	
	SYSLOCS				
	VMBLOK				
DMKCFTRM	TERMINAL command processor.		1		DMKCFM
DMKCKP	Save pertinent data when check point occurs			DMKSAVRS	
	<u>Data Areas Used:</u>				
	RCHBLOK				
	RCUBLOK				
	RDEVBLCK				
	SAVEAREA				
	VMBLOK				
DMKCKPT	Check Point program.		1		DMKSAV
DMKCNS	Real Console Terminal Manager				
DMKCNSD			20		
DMKCNSID			9		
DMKCNSIN			1		DMKIOS
DMKCNSNM			15		
DMKCNSOF			11		DMKCPV DMKIOS DMKQCN
DMKCPB	Simulate the operator's console for the virtual machine.			DMKCFPRD DMKCFPRR DMKCVTBH DMKCVTHB DMKERMSG DMKFREE DMKFRET DMKIOSQR DMKPGSPO DMKPTRAN DMKQCNWT DMKSCNFD DMKSCNVU DMKVATBC DMKVATMD	
	<u>Data Areas Used:</u>				
	IOBLOK				
	RDEVBLK				
	SAVEAREA				
	VCHBLOK				
	VCUBLOK				
	VDEVBLK				
	VMBLOK				
DMKCPBEX	Process the EXTERNAL command.		2		DMKCFM
DMKCPBNR	Process the NOTREADY command.		3		DMKCFM
DMKCPBRS	Process the RESET command.		3		DMKCFM
DMKCPBRW	Process the REWIND command.		4		DMKCFM
DMKCPBRY	Process the READY command.		3		DMKCFM
DMKCPBSR	Process the SYSTEM command.		1		DMKCFM
DMKCPI	Prepare VM/370 for operation			DMKBLDRT	

Name	Comments	MOD	Chart	Calls To	Called By
DMKCPI (cont)	<u>Data Areas Used:</u> ALOCBLOK CORTABLE PAGTABLE RDEVBLK SEGTABLE SWPTABLE VMBLOK			DMKCFMEN DMKCOGFI DMKCSOSD DMKCVTBD DMKCVTBH DMKCVTDT DMKDSPCH DMKPRETR DMKIOEFL DMKIOSQR DMKLOGOP DMKQCNRD DMKSCHST DMKSCNRD DMKSCNVS DMKWRMST	
DMKCIEM	Saves necessary data for automatic re-IPL.		2		DMKCNS
DMKCPINT	Start initialization of the VM/370 control program.	5B1	1		DMKSAV
DMKCPV	Command Processor <u>Data Areas Used:</u> BUFFER CORTABLE IOBLOK RDEVBLK SAVEAREA			DMKACODV DMKACOFF DMKACOTM DMKCVTHB DMKCVTBD DMKCVTBH DMKDMPRS DMKERMSG DMKFREE DMKFRET DMKIOSQR DMKPTRAN DMKPTRUL DMKQCNRD DMKQCNWT DMKRPAPT DMKSCNAU DMKSCNFD DMKSCNRU DMKSCNVS DMKSCNVU	

Name	Comments	MOD	Chart	Calls To	Called By
DMKCPVAC	ACNT command processor.			6	DMKCFM
DMKCPVDS	DISABLE command processor.			3	DMKCFM
DMKCPVEN	ENABLE command processor.			1	DMKCFM
DMKCPVH	HALT command processor.			12	DMKCFM
DMKCPVLK	LOCK command processor.			4	DMKCFM
DMKCPVRY	VARY command processor.			7	DMKCFM
DMKCPVSH	SHUTDOWN command processor.			5	DMKCFM
DMKCPVUL	UNLOCK command processor.			5	DMKCFM
DMKCQG	Process QUERY command. <u>Data Areas Used:</u> BUFFER SAVEAREA SFBLOK VCHBLOK VCUBLOK VDEVBLOK VMBLOK				DMKACOTM DMKCVTBD DMKCVTBH DMKCVTDB DMKCVTDT DMKCVTHB DMKERMSG DMKFREE DMKFRET DMKQCNWT DMKSCNAU DMKSCNFD DMKSCNRD DMKSCNVN DMKSCNVU
DMKCQGEN	QUERY command processor for Class G users.			1	DMKCFM
DMKCQFI	QUERY command processor for Class G users.			1	DMKCPI
DMKCQGLG	QUERY command processor for Class G users.			2	DMKLOG DMKLOG
DMKCQP	Process QUERY command <u>Data Areas Used:</u> RCHBLOK RCUBLOK RDEVBLOK SAVEAREA VCHBLOK VCUBLOK VDEVBLOK VMBLOK				DMKCVTBD DMKCVTBH DMKCVTHB DMKERMSG DMKFREE DMKFRET DMKQCNWT DMKSCNAU DMKSCNFD DMKSCNRD DMKSCNRN DMKSCNRU DMKSCNVD DMKSCNVU



Name	Comments	MOD	Chart	Calls To	Called By
DMKCQPRV	QUERY command processor for Classes B, E, and G users.		1		DMKCFM
DMKCSO	Process real spooling commands for real unit record devices. <u>Data Areas Used:</u> IOBLOK RDEVBLCK SAVEAREA SAVEWRK2 SFBLOK VMBLOK			DMKCVTBH DMKCVTDB DMKCVTHB DMKDSPCH DMKERMSG DMKFREE DMKFRET DMKIOSQR DMKPTRAN DMKPTRUL DMKQCNWT DMKRSPX DMKSCNFD DMKSCNRU DMKSCNVU DMKSPLDL DMKSTKIO	
DMKCSOBS	BACKSPACE command processor.		7		DMKCFM
DMKCSODR	DRAIN command processor.		5		DMKCFM
DMKCSOFL	FLUSH command processor.		1		DMKCFM
DMKCSOLD	LCADBUF command processor.		7		DMKCFM
DMKCSORP	REPEAT command processor.		6		DMKCFM
DMKCSOSD	Start entry point for warm start.		5		DMKCPI DMKCSP DMKCSU
DMKCSOSP	SPACE command processor.		2		DMKCFM
DMKCSOST	START command processor.		3		DMKCFM
DMKCSOVL	LOAD virtual Forms Control Buffer.		12		DMKCFM
DMKCSP	Process Class D and G spooling commands. <u>Data Areas Used:</u> SAVEAREA SFBLOK VDEVBLCK VMBLOK			DMKDCOSD DMKCVTDB DMKCVTHB DMKERMSG DMKFREE DMKFRET DMKQCNWT DMKSCNFD DMKSCNVU DMKSPLDC DMKUDRFU DMKVSPCO DMKVSPCR	

Name	Comments	MOD	Chart	Calls To	Called By
DMKCSPL	CLOSE command processor.		1		DMKCFM
DMKSPFR	FREE command processor.		3		DMKCFM
DMKCSPL	HOLD command processor.		3		DMKCFM
DMKSPSP	SFOCL command processor.		4		DMKCFM
DMKCSU	Process class D and G spooling commands. <u>Data Areas Used:</u> SAVEAREA SFBLOK VDEVBLK VMBLOK			DMKCSOSD DMKCVTBD DMKCVTDB DMKERMSG DMKFREE DMKFRET DMKQCNWT DMKSCNFD DMKCSNAU DMKSPLDL DMKUDRFU	
DMKCSUCH	CHANGE command processor.		1		DMKCFM
DMKCSUOR	ORDER command processor.		4		DMKCFM
DMKCSUPU	PURGE command processor.		6		DMKCFM
DMKCSUTR	TRANSFER command processor.		8		DMKCFM
DMKCVT	Convert Routines <u>Data Areas Used:</u> BALRSAVE TEMPSAVE			None	
DMKCVTBD	Convert binary to EBCDIC decimal.		3		DMKCDB DMKCPI DMKCPV DMKCQG DMKCQP DMKCSU DMKDEF DMKDIA DMKLNK DMKLOG DMKRSP DMKSEP DMKSPL DMKUSO DMKVDB
DMKCVTBH	Convert binary to EBCDIC hexadecimal.		1		DMKACO DMKBLD DMKCCH DMKCDB DMKCDS DMKCFD

Name	Comments	MOD	Chart	Calls To	Called By
					DMKCFP
					DMKCFS
					DMKCQG
					DMKCQP
					DMKCPI
					DMKCPV
					DMKCSO
					DMKDEF
					DMKDIA
					DMKLNK
					DMKLOG
					DMKMSW
					DMKPSA
					DMKRSP
					DMKSCH
					DMKSEP
					DMKSPL
					DMKTRC
					DMKUSO
					DMKVCA
					DMKVDB
DMKCVTDB	Convert EBCDIC decimal to binary.		3		DMKCDB
					DMKCDS
					DMKCFP
					DMKCFS
					DMKCFT
					DMKCQG
					DMKCSO
					DMKCSP
					DMKCSU
					DMKDEF
					DMKMSG
DMKCVTDT	Convert data and time to EBCDIC.		4		DMKCFS
					DMKCPI
					DMKCQG
					DMKHVC
					DMKIOF
					DMKLOG
					DMKMID
					DMKQCN
					DMKRSP
					DMKSEP
					DMKSPL
					DMKUSO
					DMKVSP

Name	Comments	MOD	Chart	Calls To	Called By
DMKCVTFP	Convert floating-point hexadecimal to binary.		1		DMKCDB
DMKCVTHB	Convert EBCDIC hexadecimal to binary.		1		DMKCDB DMKCDS DMKCFD DMKCFM DMKCFP DMKCFS DMKCPV DMKCQG DMKCQP DMKCSO DMKCSP DMKDEF DMKDIA DMKLNK DMKVDB
DMKDAS	DASD ERP <u>Data Areas Used:</u> IOBLOK ICERBLCK RDEVBLOK SAVEAREA VMBLOK			DMKCVTBH DMKFREE DMKFRET DMKIOESD DMKIOSQR DMKMSWR DMKQCNWT DMKSCNRU	
DMKDASER	Retry the failing DASD channel program.		1		DMKIOS
DMKDASRD	Process unsolicited Device End interrupts.		9		DMKIOS
DMKDASSD	Collect 3330 Statistical Data		8		DMKCPV
DMKDEF	Define a virtual device or storage. <u>Data Areas Used:</u> VDEVBLOK VMBLOK SAVEAREA			DMKBLDRL DMKBLDRT DMKCFPRD DMKCFPRR DMKCVTBD DMKCVTBH DMKCVTDE DMKCVTHB DMKERMSG DMKFREE DMKPGSPO DMKQCNWT DMKSCNFD	

Name	Comments	MOD	Chart	Calls To	Called By
				DMKSCNVD DMKSCNVN DMKSCNVU DMKUDRFU DMKUDRLK DMKUDRRD DMKUDRRV DMKUDRUL DMKVDSDF	
DMKDEFIN	DEFINE command processor.		1		DMKCFM
DMKDGD	DASD I/O <u>Data Areas Used:</u> IOBLOK SAVEAREA VDEVBLCK VMBLOK			DMKFREE DMKFRET DMKIOSQV DMKPTRUL DMKSCNVU	
DMKDGDDK	Perform disk I/O.		1		DMKHVC
DMKDIA	Connect terminal to virtual 270X or convert virtual channel to channel adapters. <u>Data Areas Used:</u> SAVEAREA VCHBLOK VCUBLOK VDEVBLOK VMBLOK			DMKACODV DMKBLDVM DMKCVTBD DMKCVTEH DMKCVTHB DMKDSPCH DMKERMSG DMKFREE DMKFRET DMKIOSQR DMKQCNWT DMKSCNAU DMKCNSFD DMKSCNRD DMKSCNVD DMKSCNVU DMKSTKIO	
DMKDIACP	COUPLE command processor.		8		
DMKDIADR	Drop dialed line to virtual System		5		DMKCCW DMKCFD
DMKDIAL	DIAL command processor.		1		DMKCFM
DMKDIASM	Simulate status for undialed lines		6		DMKVCA
DMKEMP	Dump system and re-IPL				
DMKDMPDK	WRITE dump to output device.		1		DMKPRG DMKPSA

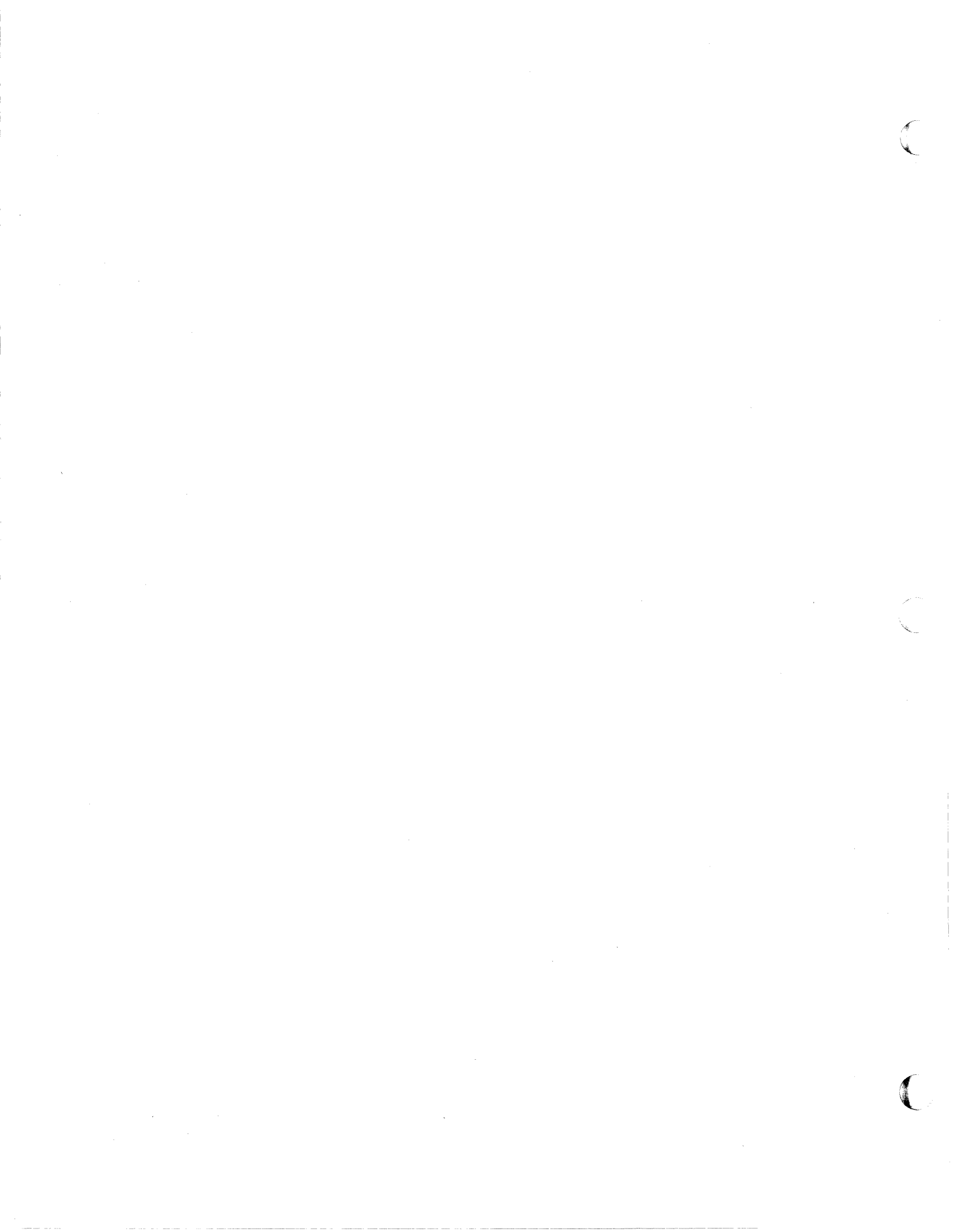
Name	Comments	MOD	Chart	Calls To	Called By
DMKDMPRS	Re-IPL system.		3		DMKCPV DMKCPV
DMKDRD	Process spool files <u>Data Areas Used:</u> SAVEAREA SFBLOK SWPTABLE UDEVELCK VSPCTL VMBLOK			DMKFREE DMKPGTSD DMKQRVKY DMKPTRAN DMKRPAGT DMKSCNVU DMKVSPCR	
DMKDRDDD	Delete system dump spool file.		7		DMKSPL
DMKDRDER	Diagnose interface to input spool files.		1		DMKHVC
DMKDRDMP	Diagnose read of system dump spool files.		6		DMKHVC
DMKDRDSY	Diagnose read of system symbol table.		7		DMKHVC
DMKDSP	Dispatcher <u>Data Areas Used:</u> CPEXBLOK ECBLOK ICBLOK VCHBLOK VCUBLOK VDEVBLOK VMBLCK			DMKCFMBK DMKQCNWT DMKSCHDL DMKSCNVU DMKTRCEX DMKTRCIO DMKTRCIT DMKUSOFF DMKVATAB DMKVATEC DMKVATMD	
DMKDSPA	Fast user re-dispatch.		9		DMKPRV
DMKDSPB	Process new virtual PSW and dispatch.		9		DMKPRG DMKPRU DMKPSA DMKVIO
DMKDSPCH	Update timers and dispatch users.	2B	1		DMKACO DMKCDB DMKCFM DMKCFD DMKCPI DMKCSO DMKDIA DMKEPS DMKGEN DMKHVC DMKIOE DMKIOS

Name	Comments	MOD	Chart	Calls To	Called By
DMKDSPCH (cont)					DMKMCH DMKPAG DMKPGT DMKPRG DMKPRV DMKPSA DMKPTR DMKRPA DMKTRC DMKTMR DMKUSO DMKVCA DMKVIO DMKVSP DMKWRM
DMKEIG	2880 Channel logout analysis. <u>Data Areas Used:</u> CCHREC				
DMKEIG80	2880 Channel logout analysis.		1		DMKCCH
DMKEPSWD	Entry point in DMKLNK.				DMKLOG
DMKERM	Message writer <u>Data Areas Used:</u> SAVEAREA VMBLOK			DMKFREE DMKFRET DMKQCNWT	
DMKERMSG	Message writer.		1		DMKACO DMKBLD DMKCDB DMKCDS DMKCFD DMKCFM DMKCFP DMKCFS DMKCFD DMKCPV DMKCQG DMKCQP DMKCSO DMKCSP DMKCSU DMKDEF DMKDIA

Name	Comments	MOD	Chart	Calls To	Called By
DMKERMSG (cont)					DMKLNK DMKLOG DMKMSG DMKRSP DMKTRA DMKUSO DMKVCH DMKVDB DMKVSP DMKWRM
DMKPRE	Free storage manager. <u>Data Areas Used:</u> FREESAVE			DMKPTRFR DMKPTRFT	
DMKFREE	Get space from free storage.	6B1	1		(General
DMKFRET	Return space to free storage.	6B2	9		entries)
DMKPRETR	Return space to free storage; do not release pages.		8		DMKCPI DMKPTR
DMKGEN	Process CMS I/O error recovery. <u>Data Areas Used:</u> IOBLOK IOERBLOK RDEVBLOK SAVEAREA VDEVBLOK VMBLOK			DMKCCWTR DMKDSPCH DMKFREE DMKFRET DMKIOSQV DMKSCNVU DMKUNTFR DMKUNTRN	
DMKGENIO	Process CMS I/O error recovery.		1		DMKHVC
DMKGRA	Read or write to the primary system console. <u>Data Areas Used:</u> None.				
DMKGRARD	Read from primary system console.		1		
DMKGRAWT	Write to primary system console.		3		DMKCKP DMKRSP
DMKHVC	Process DIAGNOSE <u>Data Areas Used:</u> RDEVBLOK VDEVBLOK VMBLOK			DMKCFMBK DMKCFMEN DMKCVTDT DMKDGDGK DMKDRDER DMKDRDMP DMKDRDSY DMKDSPCH	



Name	Comments	MOD	Chart	Calls To	Called By
DMKHVC (cont)				DMKFREE DMKGENIO DMKIOEFM DMKPGSPP DMKPRGSM DMKPRVKY DMKPTRAN DMKRPAGT DMKSCNRU DMKSCNVU DMKUDRDS	
DMKHVCAL	Process a diagnose instruction from a virtual machine.		1		DMKPRV



Name	Comments	MOD	Chart	Calls To	Called By
DMKIOE	Initiate error recording. <u>Data Areas Used:</u> IDBLOK ICERBLOK IRMBLOK RDEVBLOK SAVEAREA VMBLOK			DMKCCHRT DMKCFMBK DMKDSPCH DMKFFREE DMKFFRET DMKIOFC1 DMKIOFE1 DMKIOFIN DMKIOFM1 DMKIOFR1 DMKIOGF1 DMKIOGF2 DMKQCNWT	
DMKICECC	Channel error from SIO in DMKICS, CC=1.		5		DMKCCH
DMKIOECH	Stack channel check recording from DMKCCH.		5		DMKIOF
DMKIOECJ	Stack channel check recording from ERP.		6		DMKIOF
DMKIOEFL	Locate starting page record for recording.		6		DMKCPI
DMKIOEFM	Clear and format recording area on disk.		7		DMKHVC
DMKIOEMC	Machine check recording.		6		DMKMCH
DMKICEMH	Stack machine check recording.		6		DMKIOF
DMKIOENV	Stack environmental recording		5		DMKIOF
DMKIOEOB	Stack OBR recording.		4		DMKIOF
DMKICERC	Stack erase request.		7		DMKIOF
DMKIOERR	Schedule recording for unit check, channel data check, and hardware environmental counts.		1		DMKIOS
DMKIOESD	Record 3330 data		5		DMKDAS
DMKIOF	Perform error recording <u>Data Areas Used:</u> CPDEVCOD ICBLOK RDEVBLOK SAVEAREA VMBLOK			DMKCVTDT DMKIOECH DMKIOECJ DMKIOEMH DMKIOENV DMKIOEOB DMKIOERC DMKFFRET DMKPGTVG DMKPGTVR DMKPTRAN DMKPTRUL	

Name	Comments	MOD	Chart	Calls To	Called By
				DMKQCNWT DMKRPAGT DMKRPAPT DMKSTKCP	
DMKIOFC1	Records channel check error from SIO in DMKIOS, CC=1.		12		DMKIOE
DMKIOFE1	Records 3330 and 2305 environmental counters.		9		DMKIOE
DMKIOFIN	Initialize pointers to available recording pages.		13		DMKIOE
DMKIOFM1	Record machine checks.		13		DMKIOE
DMKIOFR1	Record outboard I/O errors.		1		DMKIOE
DMKIOG	Initialization/cleanup error recording routines. <u>Data Areas Used:</u> RDEVBLOK SAVEAREA VMBLOK			DMKFREE DMKPGTVG DMKPGTVR DMKQCNWT DMKRPAGT DMKRPAPT DMKSCNRU	
DMKIOGF1	Initialize RMS functions.		1		DMKIOE
DMKIOGF2	Erase recording area on disk.		4		DMKIOE
DMKIOS	I/O Supervisor <u>Data Areas Used:</u> CPEXBLOK IOBLOK RCHBLOK RCUBLOK RDEVBLOK SAVEAREA TEMPSAVE VDEVBLOK VMBLOK			DMKCCHIS DMKCCHNT DMKCNSIN DMKDASER DMKDSPCH DMKFREE DMKFRET DMKIOERR DMKRSPER DMKSCHDL DMKSCNRU DMKSTKCP DMKSTKIO DMKTAPER DMKTRCSI DMKVIODC	
DMKIOSHA	Halt an active device and drain all interrupts.		1		DMKCFP DMKCPV
DMKIOSIN	Process an I/O interrupt.	1B4.1	3		
DMKIOSQR	Schedule control program generated I/O.	1B4.3	1		DMKACO DMKCFP DMKCNS

Name	Comments	MOD	Chart	Calls To	Called By
DMKIOSQR (cont)					DMKCPI DMKCPV DMKCSO DMKDAS DMKDIA DMKPAG DMKPRG DMKQCN DMKRSP DMKSEP DMKSPL DMKUDR DMKVDB DMKDGD DMKGEN DMKVIO DMKCFP DMKVDB
DMKIOSQV	Schedule virtual machine I/O	1B4.3	1		
DMKIOSRW	Process IOBLOK used for rewind.		11		
DMKISM	Process ISAM CCWs <u>Data Areas Used:</u> IOBLOK RCWTASK SAVEAREA VMBLOK			DMKFREE	
DMKISMTR	Find and modify ISAM CCWs		1		DMKCCW
DMKLNK	Link to a virtual device. <u>Data Areas Used:</u> BUFFER RDEVBLOK SAVEAREA UDEVBLOK UDIRBLOK VDEVBLOK VMBLOK			DMKCVTBD DMKCVTBH DMKCVTHB DMKEPSWD DMKERMSG DMKFREE DMKFRET DMKQCNRD DMKQCNWT DMKSCNAU DMKSCNFD DMKSCNLI DMKSCNVN DMKSCNVS DMKSCNVU DMKUDRFD DMKUDRFU	

Name	Comments	MOD	Chart	Calls To	Called By
DMKLNK (cont)				DMKUDRRV DMKVDBRL DMKVDSLK	
DMKEPSWD	Enter a password. (Part of LINK)		1		DMKLOG
DMKLNKIN	LINK command processor.		1		DMKCFM
DMKLNKSB	LINK subroutines.		9		DMKLOG
DMKLOC	Lock system resource by unique name. <u>Data Areas Used:</u> CPEXBLOK			DMKFREE DMKFRET	
DMKLOCK	Lock a name.		1		DMKLNK
DMKLOCKD	Dequeue a locked name.		1		DMKDEF DMKLNK DMKUDR DMKUSO DMKVDB
DMKLOCKQ	Queue or lock a name.		2		DMKDEF DMKUDR DMKUSO DMKVDB
DMKLOCKT	Test if a name is locked.		1		
DMKLOG	Logon a user or operator. <u>Data Areas Used:</u> BUFFER RDEVBLOK SAVEAREA UDEVBLOK UDIRBLOK UMACBLOK VCHBLOK VCUBLOK VDEVBLOK VMBLOK			DMKACON DMKBLDRT DMKCFPII DMKCGFI DMKCVTBD DMKCVTBH DMKCVTDT DMKEPSWD DMKERMSG DMKFREE DMKFRET DMKLNKSB DMKQCNWT DMKSCHCP DMKSCH80 DMKSCNAU DMKSCNFD DMKSCNRD DMKSCNRU DMKSCNVN DMKSCNVU DMKTMRCK	

Name	Comments	MOD	Chart	Calls To	Called By
DMKLOG (cont)				DMKUDRFU DMKUDRRD DMKUDRRV DMKUSOFL DMKVDSAT DMKVDSDF	
DMKLOGON	LOGON a user.		1		DMKCFM
DMKLOGOP	LOGON the operator.		9		DMKCPI
DMKMCC	Monitor Command Handler <u>Data Areas Used:</u> CORTABLE MONCOMM PSA VMBLOK			DMKERMSG DMKFREE DMKPTRFR DMKPTRFT DMKQCNWT DMKSCNFD	
DMKMCCCL	Process the MONITOR Command		1		DMKCFM
DMKMCH	Machine Check Handler <u>Data Areas Used:</u> CORTABLE MCHAREA MCRECORD PAGTABLE PSA SEGTABLE SWPTABLE VMBLOK			DMKCFMBK DMKCFPRR DMKDMPRS DMKDSPCH DMKFREE DMKIOEMC DMKPGSPO DMKPTRFT DMKQCNWT DMKSCNFD DMKSTKCP DMKUSOFF	
DMKMCHIN	Process a machine check interrupt.	8B1	1		DMKCPI
DMKMCHMS	Enable or disable soft recording from SET command	8B1	8		DMKCFM
DMKMID	Date Change			DMKCVTDT DMKQCNWT DMKSCHST	
DMKMIDNT	Change system date at midnight.		1		DMKSCH

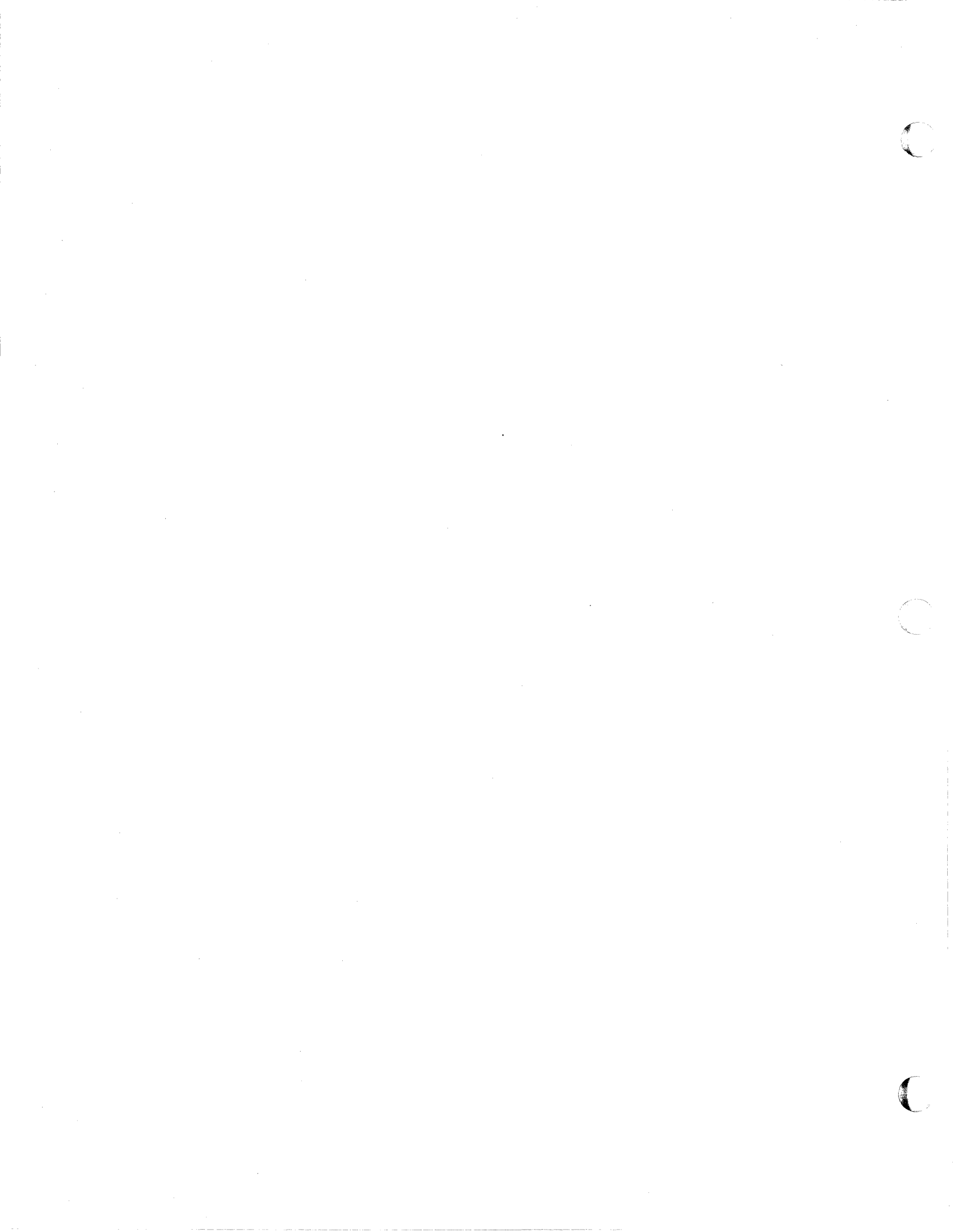
Name	Comments	MOD	Chart	Calls To	Called By
DMKMSG	Message Handler. <u>Data Areas Used:</u> BUFFER RDEVBLK SAVEAREA VDEVBLK VMBLK			DMKCVTDB DMKERMSG DMKFREE DMKFRET DMKQCNRD DMKQCNWT DMKSCNAU DMKSCNFD	
DMKMSGEC	ECHO command processor.		2		DMKCFM
DMKMSGMS	MSG command processor.		2		DMKCFM
DMKMSGWN	WNG command processor.		1		DMKCFM
DMKMSW	ERP Message Writer <u>Data Areas Used:</u> IOBLK IOERBLK RDEVBLK SAVEAREA VMBLK			DMKCVTBH DMKFREE DMKFRET DMKQCNRD DMKQCNWT DMKSCNRN	
DMKMSWR	ERP message writer.		1		DMKCNS DMKDas DMKSPL DMKTAP
DMKNEM	Translate op-codes. <u>Data Areas Used:</u> SAVEAREA				
DMKNEMOP	Translate op-codes.		1		DMKTRC
DMKPAG	Perform paging I/O <u>Data Areas Used:</u> CORTABLE CPEXBLOK IOBLK RDEVBLK SAVEAREA SWPTABLE VMBLK			DMKDSPCH DMKFREE DMKIOSQR DMKSTKCP	
DMKPAGIO	Process requests for paging I/O.	1B3.5	1		DMKPTR DMKRPA



Name	Comments	MOD	Chart	Calls To	Called By
DMKPGS	Release virtual storage. <u>Data Areas Used:</u> CORTABLE PAGTABLE SAVAREA SEGTABLE SHRTABLE SWPTABLE VMBLOK				DMKFREE DMKFRET DMKPGTPR DMKPTRAN DMKPTRFT
DMKPGSPO	Release user's entire virtual storage.	1B3.4	2		DMKCFP DMKDEF DMKMCH DMKUSO
DMKPGSPP	Release a specified area of virtual storage.	1B3.4	1		DMKHVC
DMKPGT	DASD storage management. <u>Data Areas Used:</u> ALOCBLOK CPEXBLOK PAGTABLE RDEVBLOK RECBLOK SAVEAREA SWPTABLE VMBLOK				DMKDSPCH DMKFREE DMKFRET DMKQCNWT DMKSTKCP
DMKPGTPG	Allocate DASD storage for paging.	1B3.3	1		DMKPTR
DMKPGTPR	Release DASD storage used for paging.	1B3.3	5		DMKPGS DMKPTR DMKRPA

Name	Comments	MOD	Chart	Calls To	Called By
DMKPGTPR (cont)					DMKSPL DMKVSP
DMKPGTSD	Release DASD storage used for spooling.		4		DMKDRD DMKSPL
DMKPGTSG	Allocate DASD storage for spooling.		3		DMKRSP DMKSPL DMKVSP DMKSPL
DMKPGTSR	Release DASD storage used for a complete spool file.	1B3.3	5		DMKSPL
DMKPGTVG	Allocate system virtual storage.	1B3.3	7		DMKIOF DMKIOG DMKLOG DMKRSP DMKSEP DMKSPL DMKUDR DMKVSP DMKWRM
DMKPGTVR	Release system virtual storage.		8		DMKIOF DMKIOG DMKRSP DMKSEP DMKUDR DMKVSP DMKWRM
DMKPRG	Interrupt Handler <u>Data Areas Used:</u> TEMPSAVE VMBLOK				DMKCFMBK DMKDMPDK DMKDSPB DMKDSPCH DMKPRVLG DMKPSAID DMKPTRAN DMKQCNT DMKTRCPG DMKVATPX DMKVATXS
DMKPRGIN	Hardware program interrupt.	1B3	1		DMKCPI
DMKPRGRF	Reflect SVC interrupt to virtual machine.	1B3	4		DMKPSA

Name	Comments	MOD	Chart	Calls To	Called By
DMKPRGSM	Simulate virtual program interrupt.		4		DMKHVC DMKPRV DMKTMR DMKVAT
DMKPRV	Simulate privileged operations. <u>Data Areas Used:</u> ECBLOK VMBLOK			DMKDS PA DMKDSPB DMKDSPCH DMKHVCAL DMKPRGSM DMKPTRAN DMKTMRN DMKTRCPV DMKVATAB DMKVATEX DMKVATLA DMKVATRN DMKVIOEX	
DMKPRVKY	Process virtual storage keys.		4		DMKDRD DMKHVC DMKTMR DMKTRC DMKPRG
DMKPRVLG DMKPSA	Simulate a privileged operation. Interrupt Handler <u>Data Areas Used:</u> SAVEAREA SYSLOCS	1B3.7	1	DMKCFMBK DMKCVTBH DMKDMPDK DMKDSPB DMKFREE DMKPRGRF DMKPTRAN DMKPTRUL DMKQCNCL DMKSCNRD DMKSTKIO DMKTRCSV	
DMKPSADU	Force an SVC 0 type of dump.		2		DMKCPI
DMKPSAEX	External interrupt handler.	1B2	4		DMKCPI
DMKPSAID	Get virtual address for any instruction.		3		DMKPRG



Name	Comments	MOD	Chart	Calls To	Called By
DMKPSARR	Get virtual address for RR instruction.		3		DMKTRC
DMKPSARS	Get virtual address for RS,SI, or SS instruction.		3		DMKTRC
DMKPSARX	Get virtual address for RX instruction.		3		DMKTRC
DMKPSASV	SVC interrupt handler.	1B1	1		DMKCPI
DMKPTR	Real Storage Manager <u>Data Areas Used:</u> CORTABLE PAGTABLE SAVEAREA SEGTABLE SWPTABLE VMBLOK			DMKCFMBK DMKDSPCH DMKFREE DMKFRET DMKFRETR DMKPAGIO DMKPGTPG DMKPGTPR DMKPTRFT DMKQCNWT DMKSCHDL DMKSTKCP	
DMKPTRAN	Translate user virtual storage address to a real storage address.	1B3.2	1		DMKCCW DMKCDB DMKCDS DMKCPV DMKCSO DMKDRD DMKHVC DMKIOF DMKPGS DMKPRG DMKPRV DMKPSA DMKRPA DMKSPL DMKTMR DMKTRC DMKVCA DMKVIO DMKVSP
DMKPTRFD	Release page after writing		11		
DMKPTRFE	Page must be swapped to extend		7		
DMKPTRFR	Get a page of real storage.	1B3.2	6		DMKCPI DMKFRE

Name	Comments	Code	Length	Called By
DMKPTRFL	Release a page of real storage.	1B3.2	1	DMKTRPT DMKTRER DMKTRCF DMKTRML DMKTRCM DMKTRSA DMKTRSL DMKTRCO DMKTRCP DMKTRSO DMKTRGD DMKTRIO DMKTRSA DMKTRPA DMKTRSE DMKTRSL DMKTRUN DMKTRCA DMKTRSP
DMKPTRLK	Lock a page of real storage.		13	
DMKPTRUL	Unlock a page of real storage.	1B3.2	12	
DMKQCN	Console queue manager <u>Data Areas Used</u> BUFFER CFEXBLOK CCNTASK IOBLOK RDEVELCK SAVEAREA VMBLOK			DMKCVTDT DMKFREE DMKFRET DMKIOSQR DMKSTKCP DMKUSOFF
DMKQCNCL	Clear message queue.		3	DMKCFS DMKPSA
DMKQCNRD	Queue console read request.		1	DMKCFM DMKCFS DMKCPI DMKCPV DMKEPS DMKLNK DMKMSG DMKMSW
DMKQCNWT	Queue console write request.		1	DMKACO DMKCCH DMKCDB DMKCDS DMKCFD

Name	Comments	MOD	Chart	Calls To	Called By
DMKQCNWT (cont)					DMKCFM DMKCFP DMKCFS DMKCFT DMKCPI DMKCPV DMKCQG DMKCQP DMKCSO DMKCSP DMKCSU DMKDEF DMKDIA DMKDSP DMKEPS DMKERM DMKIOE DMKIOF DMKIOG DMKLNK DMKLOG DMKMCH DMKMSG DMKMSW DMKPGT DMKPRG DMKPSA DMKPTR DMKRSP DMKSPL DMKTRA DMKTRC DMKUSO DMKVCA DMKVCH DMKVCN DMKVDB DMKWRM
DMKRPA	Virtual storage mapping. <u>Data Areas Used:</u> CORTABLE PAGTABLE SAVEAREA SWPTABLE VMBLOK				DMKDSPCH DMKFREE DMKPAGIO DMKPGTPR DMKPTRAN DMKPTRFT DMKPTRUL DMKSCHDL

Name	Comments	MOD	Chart	Calls To	Called By
DMKRPAGT	Page-in from DASD to user's virtual storage.		1		DMKCFP DMKDRD DMKHVC DMKIOF DMKIOG DMKRSP DMKUDR DMKVSP DMKWARM
DMKRPAPT	Page out to DASD from User's virtual storage.		1		DMKCVP DMKIOF DMKIOG DMKRSP DMKVSP DMKWARM
DMKRSE	Real U/R device I/O error handler. <u>Data Areas Used:</u> IOBLOK IOERBLOK RDEVBLOK SAVEAREA VMBLOK			DMKFRET DMKMSWR	
DMKRSERR	Real spool error processing.		1		DMKRSP
DMKRSP	Real spooling manager <u>Data Areas Used:</u> IOBLOK RDEVBLOK RSPLCTL SFBLOK VMBLOK			DMKACOPU DMKCVTBD DMKCVTBH DMKCVTDT DMKERMSG DMKFREE DMKFRET DMKIOSQR DMKPGTSG DMKPGTVG DMKPGTVR DMKQCNWT DMKRPAGT DMKRPAPT DMKRSERR DMKSCNFD DMKSCNRU DMKSEPS	



Name	Comments	MOD	Chart	Calls To	Called By
DMKRSP (cont)				DMKSPLCR DMKSPLDL DMKSPLOR DMKUDRFU	
DMKRSPER	Processing spooling errors (ERP).		8		DMKIOS
DMKRSPEX	Process spooling operations for real UR devices.	4B2	1		DMKACO DMKCSO DMKIOS DMKSPL
DMKSAV	Save CP nucleus or SYSRES			DMKCPINT	



Name	Comments	MOD	Chart	Calls To	Called By
DMKSAVNC	Write a page image of the control program's nucleus onto DASD.	5B3	2		
DMKSAVRS	Restore a page image copy of the control program's nucleus from DASD into main storage.	5B3	2		DMKCKP
DMKSCH	Scheduler <u>Data Areas Used:</u> CORTABLE TEMPSAVE TRQBLOK VMBLOK			DMKCVTBH DMKFRET DMKMIDNT	
DMKSCHAE	Interrupt from expiration of execution interval.		8		DMKCFS
DMKSCHCP	Interrupt from real CPU timer.		7		DMKLOG
DMKSCHDL	Alters a user's dispatching states.	3B	1		DMKDSP DMKIOS DMKPTR DMKRPA DMKUSO DMKCPI
DMKSCHMD	Interrupt from midnight date change.		8		
DMKSCHRT	Reset a clock comparator interrupt.		7		DMKCFP DMKCFS DMKMID DMKTMR DMKUSO
DMKSCHST	Establish a clock comparator interrupt.		7		DMKCPI DMKMID DMKTMR
DMKSCH80	Interrupt from real timer at storage address 80.		6		DMKCFS DMKLOG
DMKSCN	Scan Routine <u>Data Areas Used:</u> BALRSAVE BUFFER RCHBLOK RCUBLOK VCHBLOK VCUBLOK VDEVBLOK VMBLOK			None	

Name	Comments	MOD	Chart	Calls To	Called By
DMKSCNAU	Find the specified VMBLOK.		3		DMKCFD DMKCFS DMKCPV DMKCQG DMKCQP DMKCSU DMKDIA DMKLNK DMKLOG DMKMSG DMKSPL DMKUSO DMKVDB DMKVSP
DMKSCNFD	Find next field in input buffer.		3		DMKCDB DMKCDS DMKCFD DMKCFM DMKCFP DMKCFS DMKCFT DMKCPV DMKCQG DMKCQP DMKCSO DMKCSP DMKCSU DMKDEF DMKDIA DMKEPS DMKLNK DMKLOG DMKMSG DMKRSP DMKTRC DMKUSO DMKVDB
DMKSCNLI	Find all links of a minidisk.		3		DMKLNK
DMKSCNRD	Determine real device address.		2		DMKBLD DMKCPI DMKCQG DMKCQP

Name	Comments	MOD	Chart	Calls To	Called By
DMKSCNRN	Find device name for a given real device address.		4		DMKDIA DMKLOG DMKPSA DMKTRA DMKUSO DMKVDB DMKCFD DMKCPS DMKCNS DMKCPV DMKCQP DMKCSO DMKHVC DMKIOG DMKIOS DMKLOG DMKMSW DMKRSP DMKTRC DMKVDB DMKVCH DMKWRM
DMKSCNRU	Find blocks for a specified real device.		1		DMKCCH DMKCFD DMKCPS DMKCPV DMKCQP DMKCSO DMKHVC DMKIOG DMKIOS DMKLOG DMKRSP DMKVCH DMKIDE DMKWRM
DMKSCNVD	Determine virtual device address.				DMKCFD DMKCPV DMKCQP DMKCSO DMKHVC DMKIOG DMKIOS DMKLOG DMKRSP DMKVCH DMKIDE DMKWRM

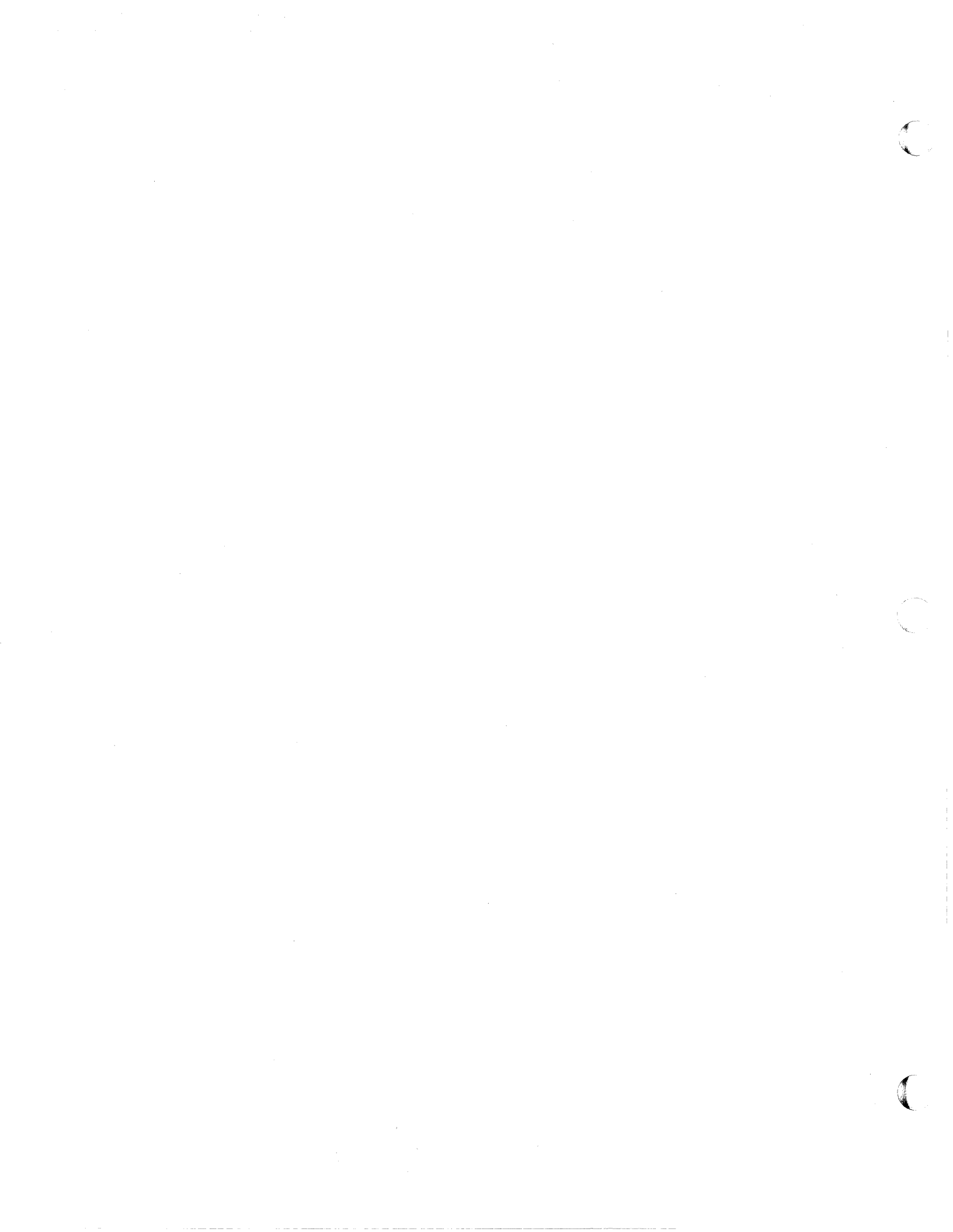
Name	Comments	MOD	Chart	Calls To	Called By
DMKSCNVN	Find device name for a given virtual device address.		4		DMKCQG DMKDEF DMKLNK DMKLOG DMKTRC DMKVDB
DMKSCNVS	Find device for specified serial number.		2		DMKCFP DMKCPI DMKCPV DMKLNK DMKLOG DMKVDB
DMKSCNVU	Find blocks for a specified virtual device address.		1		DMKCFD DMKCFM DMKCFP DMKCPV DMKCQG DMKCQP DMKCSO DMKCSP DMKDEF DMKDGD DMKDIA DMKDRD DMKDSP DMKGEN DMKHVC DMKLNK DMKLOG DMKTRC DMKVCA DMKVCH DMKVDB DMKVDG DMKVIO DMKVSP
DMKSEP	Print/punch output separator <u>Data Areas Used:</u> IOBLOK DMKEOX RDEVBLOK SAVEAREA SFBLOK				DMKCVTBD DMKCVTEH DMKCVTDT DMKDSPCH DMKIOSQR DMKPGTVG DMKPGTVR DMKPTRUL

Name	Comments	MOD	Chart	Calls To	Called By
DMKSEPSP	Print and punch the respective separators on real spooling devices.		1		DMKRSP
DMKSEV	2870 Channel logout analysis. <u>Data Areas Used:</u> CCHREC				
DMKSEV70	2870 Channel logout analysis.		1		DMKCCH
DMKSIX	2860 Channel logout analysis. <u>Data Areas Used:</u> CCHREC				
DMKSIX60	2860 Channel logout analysis.		1		DMKCCH
DMKSPL	Spool file manager <u>Data Areas Used:</u> CPEXBLOK IOBLOK IOERBLOK OWNDLIST RDEVBLOK RECBLOK RSPLCTL SAVEAREA SFBLOK VCONCTL VCUBLOK VDEVBLOK VMBLOK VSPLCTL			DMKCVTBD DMKCVTDT DMKDRDDD DMKDSPCH DMKFREE DMKFRET DMKIOSQR DMKMSWR DMKPGTSD DMKPGTSG DMKPGTSR DMKPGTVG DMKPTRAN DMKPTRLK DMKPTRUL DMKQCNWT DMKSCNAU DMKSTKCP DMKSTKIO	
DMKSPLCR	Close real reader file.		2		DMKRSP
DMKSPLCV	Close virtual printer or punch file.		1		DMKVSP
DMKSPLDL	Delete spool file buffers.		3		DMKCSO DMKCSU DMKRSP DMKVSP
DMKSPLDR	Delete spool file blocks.		3		
DMKSPLOR	Open real reader file.		1		DMKRSP

Name	Comments	MOD	Chart	Calls To	Called By
DMKSPLOV	Open virtual printer or punch file.		1		DMKVSP
DMKSSP	System Initialization/ configuration. <u>Data Areas Used:</u> RCHBLOK RCUBLOK RDEVBLOK			DMKCPINT DMKCVTBH DMKCVTHB DMKSCNRU	
DMKSSP01	Build real I/O blocks for minimum system.		1		
DMKSTK	Stack I/O <u>Data Areas Used:</u> CPEXBLOK IOBLOK				
DMKSTKCP	Stack a CPEXBLOK		1		DMKACO DMKCFM DMKCNS DMKIOF DMKIOS DMKMCH DMKPAG DMKPGT DMKPTR DMKQCN DMKSPL DMKTRC DMKUSO DMKVCA
DMKSTKIO	Stack an IOBLOK.		1		DMKACO DMKCCW DMKCFP DMKCFS DMKCSO DMKIOS DMKPSA DMKSPL DMKTMR DMKVIO
DMKTAP	Tape ERP <u>Data Areas Used:</u> IOBLOK IOERBLOK			DMKFREE DMKFRET DMKMSW	



Name	Comments	MOD	Chart	Calls To	Called By
DMKTAP (cont)	RDEVBLOK SAVEAREA VMBLOK				
DMKTAPER	Magnetic tape ERP.		1		DMKIOS
DMKTDK	T-disk space manager <u>Data Areas Used:</u> ALOCBLOK RDEVBLOK SAVEAREA				
DMKTDKGT	Get T disk space.		1		DMKVDS
DMKTDKRL	Release T disk space.		2		DMKVDE



Name	Comments	MOD	Chart	Calls To	Called By
DMKTMR	Simulate CPU timer and TOD clock <u>Data Areas Used:</u> ECBLOK TRQBLOK VMBLOK			DMKDSPCH DMKPRGSM DMKPRVKY DMKPTRAN DMKSCHRT DMKSCHST DMKSTKIO	
DMKTMRCK	Simulate virtual clock comparator interrupt.		1		DMKLOG
DMKTMRTN	Simulate timer instruction.		1		DMKPRV
DMKTRA	TRACE command processor. <u>Data Areas Used:</u> TREXT			DMKERMSG DMKFREE DMKFRET DMKQCNWT DMKSCNFD DMKTRCIT DMKTRCPB	
DMKTRACE	TRACE command processor.	7B3.2	1		DMKCFM
DMKTRC	TRACE command routines. <u>Data Areas Used:</u> TREXT			DMKCCWSB DMKCFMBK DMKCVTBH DMKDSPCH DMKFREE DMKFRET DMKNEMOP DMKPRVKY DMKPSARR DMKPSARS DMKPSARX DMKPTRAN DMKQCNWT DMKSCNRD DMKSCNRN DMKSCNVN DMKSCNVU DMKSTKCP DMKVATR DMKVSPRT	

Name	Comments	MOD	Chart	Calls To	Called By
DMKTRCEX	Process an external interrupt.	7B3.2	1		DMKDSP
DMKTRCIO	Process an I/O interrupt.	7B3.2	6		DMKDSP
DMKTRCIT	Set "SVC B2" for instruction tracing.		9		DMKDSP DMKTRA
DMKTRCND	Forcibly end tracing.		11		
DMKTRCPB	Put back user instructions altered by tracing.		9		DMKCFM DMKTRA
DMKTRCPG	Process a program interrupt.	7B3.2	6		DMKPRG
DMKTRCPV	Process a privileged instruction interrupt.	7B3.2	8		DMKPRV
DMKTRCSI	Process an I/O operation (SIO, TIO, HIO, TCH).	7B3.2	6		DMKIOS DMKVIO
DMKTRCSV	Process an SVC, Branch, or full instruction TRACE.	7B3.2	7		DMKPSA
DMKTRCSW	TRACE virtual and real CSWs.		8		DMKVIO
DMKTRM	Identify type of terminal <u>Data Areas Used:</u> RDEVBLK SAVEAREA VMBLOK				
DMKTRMID	Determine type of 2741.		1		DMKCNS
DMKUDR	Directory Manager <u>Data Areas Used:</u> IOBLOK OWNDLIST RDEVBLK SAVEAREA SYSLOCS UDEVBLK UDLBLOK UDIRBLOK UMACBLOK VMBLOK			DMKFREE DMKFRET DMKIOSQR DMKPGTVG DMKPGTVR DMKPTRAN DMKPTRUL DMKRPAGT	
DMKUDRBV	Build a list of virtual page buffers for each UDIRBLOK page on disk.		4		DMKCPI
DMKUDRDS	Swap active user directory to newly created user directory.		5		DMKHVC

Name	Comments	MOD	Chart	Calls To	Called By
DMKUDRFD	Put specified UDEVBLK into user's buffer.		3		DMKLNK DMKLOG
DMKUDRFU	Put UDIRBLOK into user's buffer.		2		DMKCSP DMKCSU DMKDEF DMKLNK DMKLOG
DMKUDRLK	Lock user directory.		1		DMKRSP DMKDEF DMKLNK DMKLOG DMKVDB
DMKUDRLT	Test user directory lock.		1		
DMKUDRRD	Put next udevblok into user's buffer.		3		DMKDEF DMKLOG
DMKUDRRV	Release a virtual page used for a buffer.		4		DMKDEF DMKLNK DMKLOG
DMKUDRUL	Unlock user directory.		1		DMKDEF DMKLNK DMKLOG DMKVDB
DMKUNT	Untranslate CCWs and CSWs <u>Data Areas Used:</u> IOBLOK RCWTASK RDEVBLK SAVEAREA VCHBLOK VDEVBLCK VMBLOK			DMKFRET DMKPTRUL DMKSTKIO	
DMKUNTFR	Release pages and free storage used for CCW chain.		1		DMKCFP DMKGEN DMKVIO
DMKUNTIS	Untranslate ISAM CCWs		5		
DMKUNTRN	Translate a real CSW to a virtual CSW.		1		DMKGEN DMKVIO
DMKUNTRS	Relocate sense byte information.		4		DMKCCW

Name	Comments	MOD	Chart	Calls To	Called By
DMKUSO	Process user termination <u>Data Areas Used:</u> RDEVBLK SAVEAREA SYSLOCS VMBLOK			DMKACODV DMKACOFF DMKACOTM DMKBLDRL DMKCFPRR DMKCVTED DMKCVTBH DMKCVTDT DMKDSPCH DMKERMSG DMKFRET DMKPGSPO DMKQCNWT DMKSCHDL DMKSCHRT DMKSCNAU DMKSCNFD DMKSCNRD DMKSTKCP DMKVATBC	
DMKUSODS	DISCONN (disconnect) command processor.		4		DMKCFM
DMKUSOFF	Logoff a user.		4		DMKCCH DMKCDS DMKDSP DMKLOG DMKMCH DMKQCN
DMKUSOFL	FORCE command processor.		3		DMKCFM DMKLOG
DMKUSCLG	LCGCFF command processor.		1		DMKCFM
DMKVAT	Storage management for EC mode virtual machine. <u>Data Areas Used:</u> BALRSAVE ECBLOK VMBLCK			DMKDSPCH DMKFREE DMKFRET DMKPRGSM DMKPTRAN	
DMKVATAB	Maintain virtual address translation.	1B3.6	1		DMKCDB DMKCDS DMKDSP DMKPRV

Name	Comments	MOD	Chart	Calls To	Called By
DMKVATBC	Return shadow tables to free storage.		3		DMKCDB DMKCDS DMKCFP DMKDSP DMKUSO
DMKVATEX	Simulate page or segment exception.		10		DMKPRV
DMKVATLA	Virtual (shadow) - virtual to virtual address translation.		5		DMKPRV
DMKVATMD	Allocate and initialization shadow table.		2		DMKCDB DMKCDS DMKCFP DMKDSP
DMKVATPX	Process paging exception for a virtual machine that performs paging.		9		DMKPRG
DMKVATRNL	Virtual (shadow) - virtual to real address translation.		7		DMKPRV DMKTRC
DMKVATSX	Process segment exception for a virtual machine that performs paging.		9		DMKPRG
DMKVCA	Simulate I/O for virtual channel to channel. <u>Data Areas Used:</u> CHXBLOK CHYBLOK IOBLOK SAVEAREA VCHBLOK VCUBLOK VDEVBLOK VMBLOK			DMKCVTBH DMKDIASM DMKDSPCH DMKFREE DMKFRET DMKPTRAN DMKPTRUL DMKQCNWT DMKSCNVU DMKSTKCP	
DMKVCARD	Reset device without de-coupling.		9		DMKCFP
DMKVCARS	Reset device and drop CTCA.		9		DMKVDB
DMKVCASH	Simulate virtual HIO.		8		DMKVIO
DMKVCAST	Simulate virtual SIO.		1		DMKVIO
DMKVCATS	Simulate virtual TIO.		8		DMKVIO
DMKVCH	Process real I/O corrections <u>Data Areas Used:</u> RCHBLOK RCUBLOK RDEVBLOK			DMKFREE DMKFRET DMKQCNWT DMKSCNRU DMKSCNVU	

Name	Comments	MOD	Chart	Calls To	Called By
	SAVEAREA			DMKVDBRL	
	VCHBLOK			DMKVDSDF	
	VCUBLOK				
	VDEVBLOK				
	VMBLOK				
DMKVCHDC	ATTACH and DETACH (real devices and channels) command processor.		1		DMKVDB
DMKVCN	Simulate user SIOs to virtual console. <u>Data Areas Used:</u>			DMKCFMEK	
	VCHBLOK			DMKDSPCH	
	VCONCTL			DMKFREE	
	VCUBLOK			DMKFRET	
	VDEVBLCK			DMKQCNRD	
	VMBLOK			DMKQCNWT	
DMKVCNEX	Simulate all SIOs to a virtual console.	7B1	1		DMKVIO
DMKVDB	Process virtual I/O connections <u>Data Areas Used:</u>			DMKACODV	
	BUFFER			DMKCFPRD	
	IOBLCK			DMKCVTBD	
	RDEVBLOK			DMKCVTBH	
	SAVEAREA			DMKCVTHB	
	VCHBLOK			DMKERMSG	
	VCUBLOK			DMKFREE	
	VDEVBLOK			DMKFRET	
	VMBLCK			DMKIOSQR	
				DMKQCNWT	
				DMKSCNAU	
				DMKSCNFD	
				DMKSCNRD	
				DMKSCNRN	
				DMKSCNRU	
				DMKSCNVN	
				DMKSCNVS	
				DMKSCNVU	
				DMKTDKRL	
				DMKUDRLK	
				DMKUDRUL	
				DMKVCARS	
				DMKVCHDC	
				DMKVDBRL	
				DMKVDSAT	
				DMKVSPCO	
				DMKVSPCR	



Name	Comments	MOD	Chart	Calls To	Called By
DMKVDBAT	ATTACH (virtual devices or channels to a virtual machine) command processor.		1		DMKCFM
DMKVDBDE	DETACH (virtual devices or channels from virtual machine) command processor.		13		DMKCFM
DMKVDBRL	Release a device from a virtual machine.		14		DMKCFP DMKLNK DMKVCH
DMKVDS	Virtual device interface <u>Data Area Used:</u> RDEVBLOK SAVEAREA VDEVBLOK VCHBLOK VCUBLOK VDEVBLOK VMBLOK			DMKFREE DMKFRET DMKSCNVU DMKTDKGT	
DMKVDSAT	Attach a virtual device.		1		DMKLOG DMKVDB
DMKVDSDF	Define a virtual device.		2		DMKDEF DMKLOG DMKVCH
DMKVDSLK	Link to a virtual DASD device.		3		DMKLNK DMKLOG
DMKVIO	Virtual I/O manager <u>Data Areas Used:</u> IOBLOK VCHBLOK VCUBLOK VDEVBLOK VMBLCK			DMKCCWTR DMKDSPB DMKDSPCH DMKFREE DMKFRET DMKIOSQV DMKPTRAN DMKSCNVU DMKSTKIO DMKTRCSI DMKTRCSW DMKUNTRN DMKUNTRN DMKVCASH DMKVCAST DMKVCATS DMKVCNEX DMKVSPEX	
DMKVIODC	Dedicated Channel Interrupt		13		DMKIOS

Name	Comments	MOD	Chart	Calls To	Called By
DMKVIOEX	Simulate a SIO, TIO, HIO, or TCH.	1B3.9	1		DMKPRV
DMKVIOIN	Translate a virtual I/O interrupt.	1B4.2	9		DMKIOS DMKSPL
DMKVMI	IPL virtual machine			None	
DMKVMIPL	Simulate IPL to a virtual machine.		1		DMKCFP
DMKVSP	Virtual spooling manager			DMKCVTBH	
	<u>Data Areas Used:</u>			DMKCVTDT	
	SAVEAREA			DMKERMMSG	
	SFBLOK			DMKFREE	
	VCHBLOK			DMKFRET	
	VCUBLOK			DMKPGTSG	
	VDEVBLOK			DMKPGTVG	
	VFCBLOK			DMKPGTVR	
	VSPLCTL			DMKPTRAN	
	VMBLOK			DMKPTRUL	
				DMKQCNWT	
				DMKRPAGT	
				DMKRPAPT	
				DMKSCNVD	
				DMKSCNVU	
				DMKSPLCV	
				DMKSPLDL	
				DMKSPLOV	
DMKVSPCO	Close spooled printers and punches.		11		DMKCDB DMKCFD DMKCSP DMKVDB
DMKVSPCR	Stop and clear all pending status for spooled card reader.		10		DMKCFP DMKCSP DMKDRD DMKTRA DMKVDB

Name	Comments	MOD	Chart	Calls To	Called By
DMKVSPEX	Simulate SIO to a spooled unit record device.		1		DMKVIO
DMKVSPRT	Put a CP generated line on the User's spooled printer		14		DMKCPB DMKTRC
DMKVSPVP	Simulate SIO to a spooled virtual console.		13		DMKVCN
DMKWRM	Warm Start			DMKDSPCH DMKERMSG DMKFREE DMKFRET DMKPGTVG DMKPGTVR DMKQCNWT DMKRPAGT DMKRPAPT DMKSCNRU	
DMKWRMST	Warm start.	5B1.4	1		DMKCPI

SUBROUTINE DIRECTORY

<u>Subroutine</u>	<u>Module</u>	<u>Chart</u>	<u>Subroutine</u>	<u>Module</u>	<u>Chart</u>
ACNTDED	DMKCPV	6	CNTRLNBW	DMKCCW	13
ACTPIRA	DMKACO	3	CNTRLSUB	DMKCCW	9
ACTPUNCH	DMKACO	3	CNVTBIN	DMKCDB	11
ADDCHEK	DMKDRD	4	CONFRRTN	DMKCFM	3
ADDCHEK	DMKHVC	2	CONNECT	DMKDRD	5
ADDQ	DMKSCH	3	CONVERT	DMKMSW	3
ADDRTEST	DMKVMI	3	COPYIOB	DMKIOS	10
ADJTIME	DMKMID	2	COPYSEGT	DMKVAT	3
ADTRANS	DMKPRV	5	CORTBLR	DMKMCH	6
ARUNLST	DMKSCH	5	CPIPINT	DMKCPI	1
ASTERISK	DMKCFM	3	CUSCAN	DMKVIO	8
AWAITLST	DMKSCH	5	CVTRADD	DMKVDB	11
BLANKBUF	DMKTRC	8	DASDXA	DMKCCW	19
BLDBLOK	DMKVDS	4	DASDXB	DMKCCW	19
BLDVDEV	DMKVDS	3	DASDXC	DMKCCW	19
BLOKLETR	DMKSEP	4	DASDXD	DMKCCW	19
BUILDCTL	DMKSPL	2	DASDXE	DMKCCW	19
CALLERR	DMKIOS	15	DASDXF	DMKCCW	20
CCWCHKEY	DMKCCW	30	DASDX0	DMKCCW	17
CCWFETCH	DMKVCA	3	DASDX1	DMKCCW	13
CCWPUTSK	DMKCCW	20	DASDX2	DMKCCW	14
CCWRELCH	DMKCCW	21	DASDX3	DMKCCW	15
CFPIO	DMKCFP	6	DASDX4	DMKCCW	15
CHACT	DMKCKP	4	DASDX5	DMKCCW	16
CHEKPTE	DMKVAT	11	DASDX6	DMKCCW	17
CHEKRKEY	DMKDGD	8	DASDX7	DMKCCW	17
CHKBUSY	DMKDRD	6	DASDX9	DMKCCW	18
CHKFORM	DMKVAT	1	DEDDXA	DMKCCW	22
CHKRDEV	DMKVDB	10	DEDDXB	DMKCCW	22
CHNREXIT	DMKIOE	4	DEDDXC	DMKCCW	23
CHSCAN	DMKVIO	8	DEDDXD	DMKCCW	23
CKBKMSG	DMKRSP	5	DEDDXE	DMKCCW	23
CKEXTRA	DMKCFP	3	DEDDXF	DMKCCW	23
CKEXTRA	DMKCPV	11	DEDDX1	DMKCCW	22
CKIOB	DMKDAS	9	DEDDX2	DMKCCW	22
CKIOB	DMKTAP	10	DEDDX3	DMKCCW	22
CKRECOMP	DMKRPA	2	DEDDX4	DMKCCW	22
CKTIME	DMKPTR	4	DEDDX5	DMKCCW	22
CKWAIT	DMKPTR	4	DEDDX6	DMKCCW	22
CLOCKFMT	DMKDMP	4	DEDDX7	DMKCCW	22
CLOSE	DMKCKP	5	DEDDX9	DMKCCW	22
CLRCHAN	DMKDMP	8	DELETE	DMKRSP	4
CMNEXT	DMKDRD	7	DELIRA	DMKSPL	9
			DEVCARD	DMKACO	1
			DEVCARD	DMKCKP	6
			DEVSCAN	DMKVIO	9
			DGRETURN	DMKDGD	7

<u>Subroutine</u>	<u>Module</u>	<u>Chart</u>	<u>Subroutine</u>	<u>Module</u>	<u>Chart</u>
DGTRANS	DMKDGD	6	FILECLR	DMKVSP	10
DGTRANSO	DMKDGD	6	FINDEVIC	DMKPGT	5
DIAGRTRN	DMKGEN	1	FINDREC	DMKPGT	6
DIALING	DMKDIA	4	FIOER	DMKIOE	5
DIALIRA	DMKDIA	4	FIOER	DMKIOF	6
DIALXA	DMKCCW	28	FIRSTREC	DMKCKP	5
DIALXB	DMKCCW	28	FMTREAD	DMKIOG	4
DIALXC	DMKCCW	28	FMTWRITE	DMKIOG	4
DIALXD	DMKCCW	28	FNDVADDR	DMKCFP	7
DIALXE	DMKCCW	28	FNDVADDR	DMKCPB	3, 4
DIALXF	DMKCCW	28	FRETCORE	DMKCQG	2
DIALX1	DMKCCW	26	FRETIOB	DMKCNS	4
DIALX2	DMKCCW	27	FRETIOB	DMKCSO	10
DIALX3	DMKCCW	27	FRETIOER	DMKCNS	7
DIALX4	DMKCCW	27	FRETPTB	DMKDAS	9
DIALX5	DMKCCW	27	FRETPTB	DMKTAP	10
DIALX6	DMKCCW	27	FRETVMB	DMKDIA	5
DIALX7	DMKCCW	28	FRET05	DMKFRE	3
DIALX9	DMKCCW	28	FREUSER	DMKDEF	6
DISASUB	DMKCPV	3	FREUSER	DMKVDB	16
DISCOMM	DMKCDB	2	FSINGIOE	DMKIOF	5
DISCRETN	DMKUSO	5	FUMAT	DMKMCH	6
DISDMPID	DMKCDB	9	GETADDR	DMKPRV	5
DISHEAD	DMKCDB	9	GETADDR	DMKRSP	7
DISINIT	DMKCDB	7	GETBUF	DMKQCN	3
DISNEXTA	DMKCDB	3	GETBUFF	DMKUDR	3
DISWRITE	DMKCDB	10	GETCCW	DMKVCN	3
DISWRTR	DMKCDB	10	GETCCW	DMKVSP	15
DMPMCHCK	DMKDMP	5	GETCHAIN	DMKCSP	10
DMPPRGCK	DMKDMP	5	GETCHAIN	DMKCSU	16
DRAINMSG	DMKRSP	7	GETCLASS	DMKCSP	7
DROPQ	DMKSCH	4	GETCLASS	DMKCSU	12
DRDLCSE	DMKDRD	6	GETCOPY	DMKCSP	8
DROPING	DMKDIA	5	GETCOPY	DMKCSU	13
DROPLIST	DMKSCH	5	GETDEVIC	DMKCSO	11
DSKIOINT	DMKDMP	3	GETDEVIC	DMKCSP	8
DSKWRTC1	DMKDMP	3	GETDISK	DMKWRM	4
DSKWTRC2	DMKDMP	3	GETENTRY	DMKRPA	2
DTBIN	DMKMID	1	GETFILE	DMKCSP	8
DTDEC	DMKMID	1	GETFILE	DMKCSU	13
DVICECLR	DMKVSP	11	GETID	DMKCSP	10
			GETID	DMKCSU	15
ECHORETN	DMKMSG	3	GETNAME	DMKCSP	9
ECHOWRIT	DMKMSG	3	GETNAME	DMKCSU	14
ECHOWTRT	DMKMSG	3	GETPAGT	DMKVAT	10
EMPTYCYL	DMKPGT	7	GETRCAW	DMKSPL	6
ENABTERM	DMKCPV	3	GETRCAW	DMKVIO	13
EXCLAIM	DMKCNS	13			

<u>Subroutine</u>	<u>Module</u>	<u>Chart</u>	<u>Subroutine</u>	<u>Module</u>	<u>Chart</u>
GETRDEV	DMKVDB	17	MCHGOZO	DMKMCH	9
GETRTYP	DMKVDB	17	MCHHSKP	DMKMCH	8
GETSHAD	DMKVAT	3	MOVEDATA	DMKVSP	16
GETUSER	DMKCSP	6	MSGALL	DMKMID	2
GETUSER	DMKCSU	10	MSGERR	DMKSPL	7
GETUSER	DMKSPL	3	MSGFMT	DMKMSG	2
GETVRADD	DMKPTR	5	MSGRET	DMKMCH	5
GETVSPL	DMKDRD	5	MSGSDND	DMKMSG	2
GETYPE	DMKCSP	6	MSGSUBR	DMKUSO	5
GETYPE	DMKCSU	12	MSG900E	DMKCKP	7
GTCLASSB	DMKCSP	7	MSG901E	DMKCKP	7
GTCLASSB	DMKCSU	12			
GU02	DMKCSU	10	NEXTBUFF	DMKRSP	5
			NXTFCR	DMKCNS	17
			NXTPTR	DMKDAS	7
IDAWSUB	DMKCCW	13	OPEN	DMKVSP	14
INITCON	DMKQCN	2	OPENCONT	DMKVSP	14
INITSCAN	DMKCSU	6	OTHRXA	DMKCCW	29
INSTADR	DMKPRV	4	OTHRXB	DMKCCW	29
IO	DMKVM I	1	OTHRXC	DMKCCW	29
IOEMSGWR	DMKIOG	4	OTHRXD	DMKCCW	29
IOINT	DMKSAV	1	OTHRXE	DMKCCW	29
IORETN	DMKRPA	2	OTHRXF	DMKCCW	29
IORETURN	DMKUDR	6	OTHRX1	DMKCCW	28
IOSDQCH	DMKIOS	13	OTHRX2	DMKCCW	29
IOSDQCU	DMKIOS	13	OTHRX3	DMKCCW	29
IOSDQDV	DMKIOS	12	OTHRX4	DMKCCW	29
IOSENSE	DMKIOS	14	OTHRX5	DMKCCW	29
IOSFINDP	DMKIOS	9	OTHRX6	DMKCCW	29
IOSGTIOB	DMKIOS	10	OTHRX7	DMKCCW	29
IOSIGNOR	DMKIOS	11	OTHRX9	DMKCCW	29
IOSQCH	DMKIOS	12	OVERHEAD	DMKPAG	3
IOSQCU	DMKIOS	11			
IOSQDEV	DMKIOS	11	PAGFREE	DMKPTR	10
IOSRECER	DMKIOS	14	PAGIN	DMKPTR	5
IOSTRTCH	DMKIOS	7	PARTLINE	DMKDMP	4
IOSTRTCU	DMKIOS	7	PASSRETN	DMKEPS	1
IOSTRTDV	DMKIOS	8	PCITEST	DMKVSP	16
			PERCHEK	DMKPRV	4
KEYTRAN	DMKHVC	1	PGTMSG	DMKPGT	4
			PGTMSG2	DMKMSG	4
LASTLINE	DMKDMP	4	PRGEVEN	DMKPRG	4
LBIONT	DMKSAV	3	PRI0	DMKDMP	5
LINESIZE	DMKCFT	1	PROTEST	DMKVSP	16
LINKDEF	DMKLNK	11	PROTTEST	DMKVCN	6
LINKSUB	DMKLNK	9	PRTDATA	DMKVSP	12
LOKUSER	DMKDEF	6	PRTDONE	DMKVSP	8
LOKUSER	DMKVDB	17			

<u>Subroutine</u>	<u>Module</u>	<u>Chart</u>	<u>Subroutine</u>	<u>Module</u>	<u>Chart</u>
PRTEOF	DMKVSP	7	SCAN	DMKUDR	1
PSETKEY	DMKPRV	3	SCANDEV	DMKCSO	5
PTRFDISP	DMKPTR	9	SCANRSP	DMKCQP	1
PUTBACK	DMKTRC	9	SCANSHQ	DMKCQP	1
PUTLINE	DMKDMP	4	SCPZCAW	DMKSAV	3
PUTMSG1	DMKRSP	5	SEARCHB	DMKDRD	5
PUTREC	DMKCKP	5	SELECT	DMKPTR	10
PUTSEEK	DMKCCW	20	SENSE	DMKVSP	8
			SENSMOVE	DMKVSP	13
QRYFCNT	DMKCQG	1	SEPIRA	DMKSEP	3
QRYUSRN	DMKCQG	2	SETCCW	DMKDMP	5
QRYVFMT	DMCQG	2	SETHHR	DMKSAV	1
			SETIMER	DMKSCH	7
RDLABEL	DMKVDB	11	SETOPTS	DMKVSP	7
RDLBIRA	DMKVDB	12	SETUPEX	DMKVAT	8
RDRDATA	DMKVSP	9	SFBCHAIN	DMKSPL	2
RDRID	DMKSPL	2	SFBSCAN	DMKCSU	6
READADDR	DMKSSP	3	SHADOWS	DMKVAT	2
READBUF	DMKDAS	9	SHADSET	DMKVAT	8
READLOG	DMKCF5	2	SNSRTN	DMKDAS	8
REALREAD	DMKIOF	5	SRCHPTE	DMKVAT	10
REALRETN	DMKCQP	2	STARTMSG	DMKCSO	5
REALWRT	DMKIOF	5	STOLOC	DMKCDS	2
RECGETNN	DMKCF5	3	STOSCAN	DMKCDS	1
RECHAIN	DMKRSP	6	STSTERR	DMKMCH	6
REGSPEC	DMKPRV	3	SVCFR	DMKPSA	2
RESETIMR	DMKSCH	8	SVCGET	DMKPSA	5
RESTINST	DMKCFD	4	SVCRLSE	DMKPSA	5
RESYSTEM	DMKCFP	4	SWPCALL	DMKVDB	16
RESYSTEM	DMKCPB	1	SWPUSER	DMKVDB	17
RET8	DMKWRM	5	SYSTUNLC	DMKCFG	3
RET8R1	DMKWRM	5	SYSUNLCK	DMKCFP	12
REWRITE	DMKPTR	11			
RSPMSG	DMKRSP	5	TAPERADD	DMKRSP	5
RSTMPPOFF	DMKSCH	6	TAPEXA	DMKCCW	24
RSTMPON	DMKSCH	5	TAPEXB	DMKCCW	24
RSTRIRA	DMKSEP	3	TAPEXC	DMKCCW	24
RTNBRKND	DMKCNS	20	TAPEXD	DMKCCW	25
RTNECA	DMKCNS	8	TAPEXE	DMKCCW	25
RTNEXCLM	DMKCNS	14	TAPEXF	DMKCCW	25
RTNIDENT	DMKCNS	10	TAPEX1	DMKCCW	23
RTNPREP	DMKCNS	12	TAPEX2	DMKCCW	23
RTNSDPRP	DMKCNS	11	TAPEX3	DMKCCW	23
RUNTIME	DMKDSP	8	TAPEX4	DMKCCW	23
			TAPEX5	DMKCCW	23
SAVRETN	DMKCFG	3	TAPEX6	DMKCCW	23
SAVRETN	DMKCPV	11	TAPEX7	DMKCCW	24
SCAN	DMKSSP	3	TAPEX9	DMKCCW	24

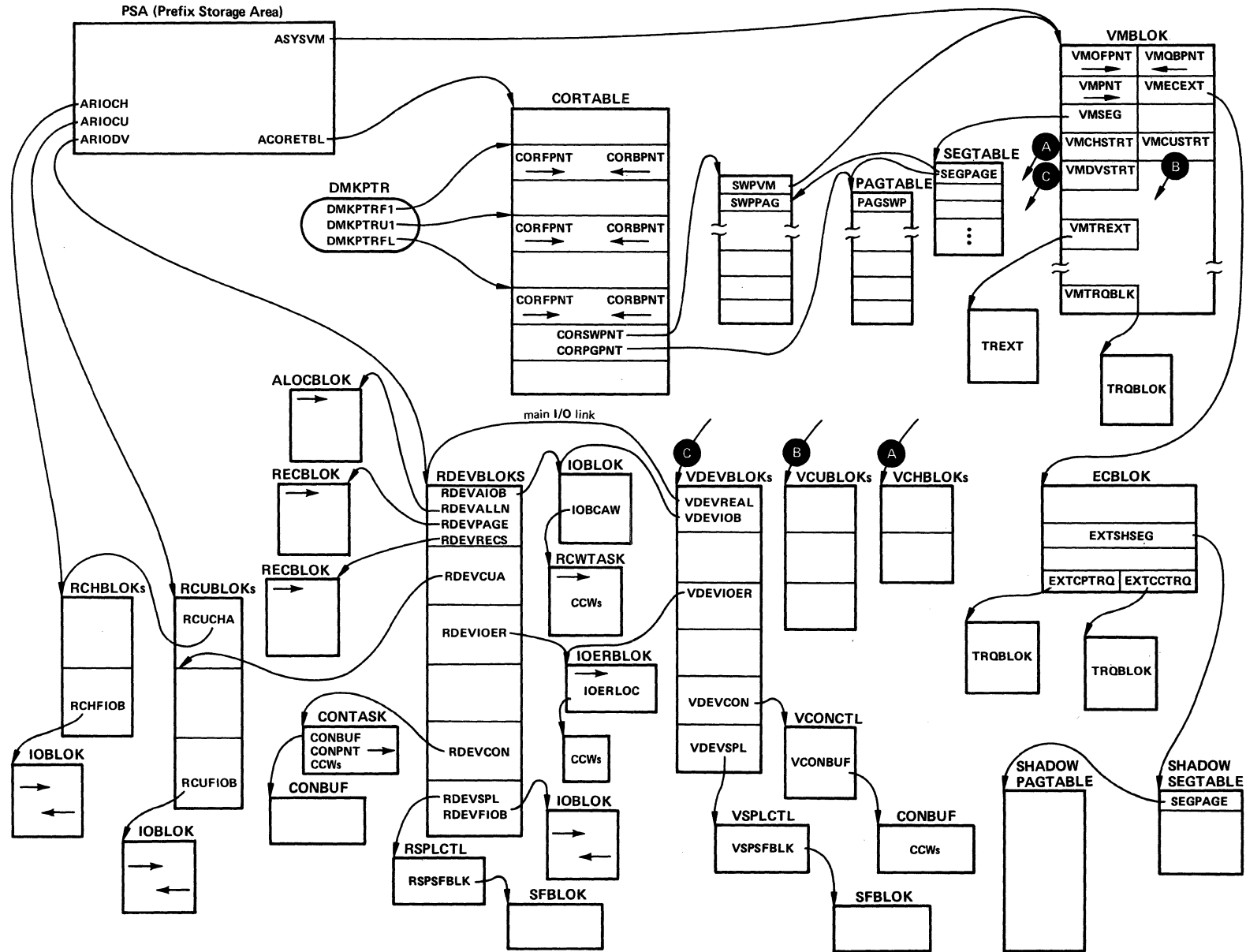
<u>Subroutine</u>	<u>Module</u>	<u>Chart</u>	<u>Subroutine</u>	<u>Module</u>	<u>Chart</u>
TDKIRA	DMKTDK	2	TSTONOFF	DMKCFS	3
TERM	DMKMCH	6	TSTONOFF	DMKCFT	2
TERMIRA	DMKRSP	6	TSTSEP	DMKSEP	3
TERMXA	DMKCCW	26	TYPLINE	DMKDMP	6
TERMXB	DMKCCW	26	UNLOKSUB	DMKLNK	11
TERMXC	DMKCCW	26	UNRELSUB	DMKUNT	3
TERMXD	DMKCCW	26	UNTRSREL	DMKUNT	5
TERMXE	DMKCCW	26	USERCARD	DMKACO	1
TERMXF	DMKCCW	26	USERCARD	DMKCKP	6
TERMX1	DMKCCW	25	USERDEV	DMKVDB	17
TERMX2	DMKCCW	25	USOSUB	DMKUSO	5
TERMX3	DMKCCW	25	VCNMV DAT	DMKVCN	4
TERMX4	DMKCCW	25	VCNRDRET	DMKVCN	8
TERMX5	DMKCCW	25	VCNRELSE	DMKVCN	6
TERMX6	DMKCCW	25	VCNSCALC	DMKVCN	4
TERMX7	DMKCCW	26	VDBSCAN	DMKVDB	1
TERMX9	DMKCCW	26	VIOINT1	DMKVIO	12
TICSUBI	DMKCCW	8	VOL1RTN	DMKDAS	9
TICSUBX	DMKCCW	8	WAITIME	DMKDSP	9
TICSUB1	DMKCCW	8	WAITPAGE	DMKPAG	2
TRAINIT	DMKTRC	8	WRTOUT	DMKCQP	1
TRANBRNG	DMKCCW	29	WRTVIRT	DMKCQG	9
TRANBRNG	DMKTRC	12	ZAPPAGE	DMKVAT	4
TRANLOCK	DMKCCW	30	ZAPSEGS	DMKVAT	4
TRANRETN	DMKPTR	4	ZAPVOLD	DMKVDB	11
TRANS	DMKCNS	19			
TROUSUB	DMKTRC	10			
TROUSUB7	DMKTRC	11			



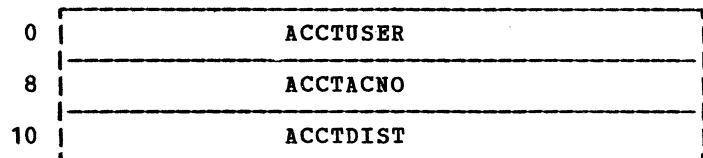
DATA AREAS -- CONTROL BLOCKS

To determine the modules that reference or alter a data area or a field in a data area, refer to the alphabetical Label Cross-reference list that is contained in the microfiche for VM/370.

Diag. 9A0. CP Control Block Relationships

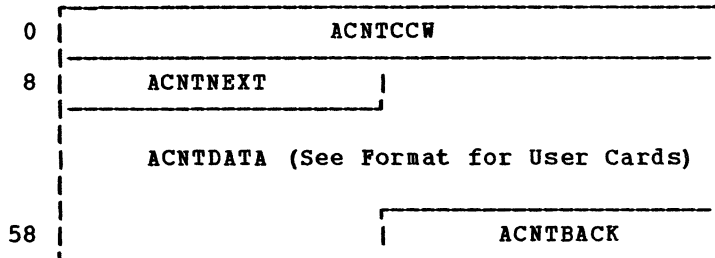


ACCTBLOK - USER ACCOUNTING BLOCK



Displacement		Field		Field Description, Contents, Meaning
Hex	Dec	Name		
0	0	ACCTUSER DS	CL8	Virtual machine identification
8	8	ACCTACNO DS	CL8	Virtual machine accounting number
10	16	ACCTDIST DS	CL8	Virtual machine distribution number
		ACCTLENG EQU	(*-ACCTBLOK)/8	Size of ACCTBLOK in doublewords (X'03')

ACNTBLOK - ACCOUNTING CARD BUFFER



Displacement		Field		Field Description, Contents, Meaning
Hex	Dec	Name		
0	0	ACNTCCW DS	D	Punch CCW for accounting card
8	8	ACNTNEXT DS	F	Address of next ACNTBLOK in chain
C	12	ACNTDATA DS	CL80	Accounting information
5C	92	ACNTBACK DS	F	Address of previous ACNTBLOK in chain
		ACNTSIZE EQU	(*-ACNTBLOK)/8	Size of ACNTBLOK in doublewords (X'0C')

Format for User Cards

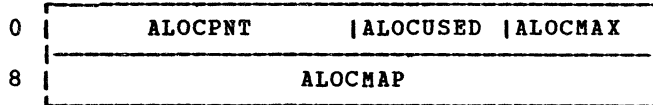
C	ACNTUSER	
14	ACNTNUM	
1C	ACNTSTOP	
		ACNTCONT
2C	ACNTTIME	ACNTVTIM
34	ACNTPGRD	ACNTPGWT
3C	ACNTIOCT	ACNTPNCH
44	ACNTLINS	ACNTPNCH
4C	ACNTRSV1	
54	ACNTRSV2	ACNTCODE

Displacement		Field	Field Description, Contents, Meaning	
Hex	Dec	Name		
		ORG	ACNTDATA	
C	12	ACNTUSER DS	CL8	Virtual machine identification
14	20	ACNTNUM DS	CL8	Virtual machine accounting number
1C	28	ACNTSTOP DS	CL12	Date and time of accounting MMDDYYHHSS
28	40	ACNTCONT DS	1F	Number of seconds connected
2C	44	ACNTTIME DS	1F	Milliseconds of CPU time used
30	48	ACNTVTIM DS	1F	Milliseconds of virtual CPU time used
34	52	ACNTPGRD DS	1F	Total page reads
38	56	ACNTPGWT DS	1F	Total page writes
3C	60	ACNTIOCT DS	1F	Virtual SIO count for non-spooled I/O
40	64	ACNTPNCH DS	1F	Virtual card count - spooled punch
44	68	ACNTLINS DS	1F	Virtual line count - spooled printer
48	72	ACNTPNCH DS	1F	Virtual card count - spooled reader
4C	76	ACNTRSV1 DS	FL8	Reserved for IBM use
54	84	ACNTRSV2 DS	HL6	Reserved for IBM use
5A	90	ACNTCODE DS	1H	Accounting card identification code
		Card code for ACNTCODE		
		DC	C'x1'	User virtual machine accounting card
		DC	C'x2'	User dedicated device accounting card
		DC	C'x3'	User temporary disk space accounting card

where x = C if the card was punched for the userid in the User Accounting Block  
 = 0 if the card was punched for the user requesting the card

2C	44	ORG	ACNTTIME	
30	48	ACNTDEVC DS	XL4	Device code (CTFM) See DEVTYPE copy file
		ACNTNCYL DS	1H	Number of cylinders of T-disk space

ALOCBLOK - DASD CYLINDER ALLOCATION BLOCK



Displacement			Field			Field Description, Contents, Meaning
Hex	Dec		Name			
0	0		ALOCPNT DS	1F		Pointer to next ALOCBLOK on chain
4	4		ALOCUSED DS	1H		Number of cylinders currently in use
6	6		ALOCMAX DS	1H		Maximum number of cylinders available
8	8		ALOCMAP DS	0F		Cylinder allocation bit map

Bits defined in ALOCMAP:  
 0 = Cylinder is available  
 1 = Cylinder has been assigned

Note: The size of the ALOCMAP is variable and depends on the number of cylinders on the device. Generally, the size of the ALOCBLOK is determined by the following formula:

$$ALOC\text{SIZE (doublewords)} = (((ALOC\text{MAX}+7)/8) + 7)/8 + 1$$

where:

ALOCMAX	for 2314	=	203
	for 3330	=	404
	for 2305-1	=	48
	for 2305-2	=	96

Any bits in the map that represent cylinders not present on the device are set to one.

For TDISK allocation blocks

		ORG	ALOCUSED	
4	4	ALOCCYL1 DS	1H	First cylinder of TDISK area
6	6	ALOCCYL2 DS	1H	Last cylinder of TDISK area

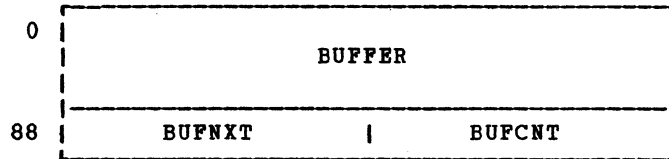
Bytes defined in ALOCMAP  
 X'00' = Cylinder is available  
 X'AA' = Cylinder has been allocated

Note: The size of the TDISK ALOCMAP is variable and depends on the number of cylinders in the range ALOCCYL1 to ALOCCYL2. Generally, the size of a given block is determined by the following formula:

$$ALOC\text{SIZE}(\text{doublewords}) = ((\text{ALOCCYL2}-\text{ALOCCYL1}+8)/8)+1$$

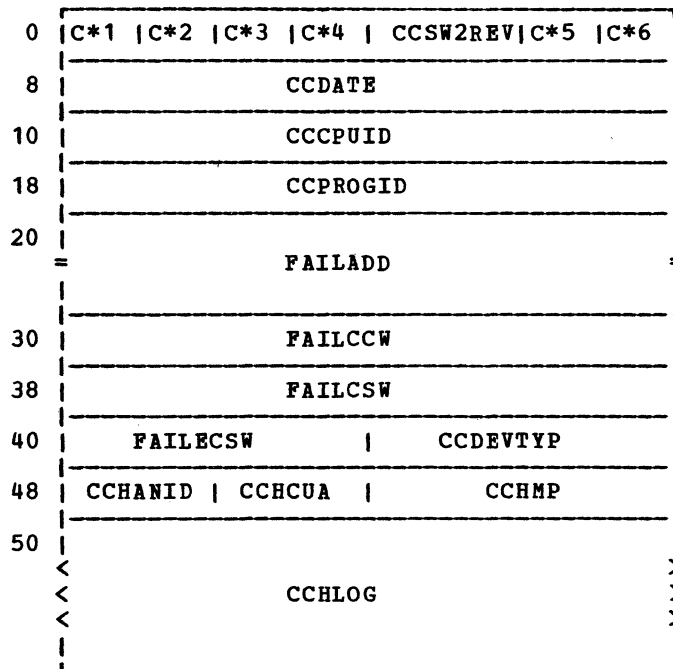
Bytes for cylinders that are not available are marked allocated.

BUFFER - CONSOLE FUNCTION INPUT BUFFER



Displacement		Field		Field Description, Contents, Meaning
Hex	Dec	Name		
0	0	BUFIN	DS CL136	Input line
88	136	BUFNXT	DS 1F	Pointer to next byte in BUFFER
8C	140	BUCFNT	DS 1F	Count of characters in input line
88	136	BUFINLTH	EQU L'BUFIN	Input BUFFER size in bytes
		BUFSIZE	EQU (*-BUFFER)/8	BUFFER size in doublewords (X'12')

CCHREC - CHANNEL CHECK HANDLER RECORD



Displacement		Field		Field Description, Contents, Meaning
Hex	Dec	Name		
0	0	CCRECTYP	DS 1X	C*1 Record type
1	1	CCOPSYS	DS 1X	C*2 Operating system
2	2	CCSW1	DS 1X	C*3 Switch one
3	3	CCSW2	DS 1X	C*4 Switch two
4	4	CCSW2REV	DS 2X	Unused
6	6	CCRECNT	DS 1X	C*5 Record count
7	7	CCRECNT1	DS 1X	C*6 Unused
8	8	CCDATE	DS 1D	Date and time
10	16	CCCPUID	DS 1D	CPUID
18	24	CCPROGID	DS 1D	USERID
20	32	FAILADD	DS 8H	Active I/O units
30	48	FAILCCW	DS 1D	Failing CCW
38	56	FAILCSW	DS 1D	Failing CSW
40	64	FAILECSW	DS 0F	Failing ECSW
40	64	IGPRGFLG	DS CL1	Program flag bits

		Bits defined in IGPRFLG	
		CCHSIOB EQU X'80'	SIO bit
		CCHINTB EQU X'40'	Interrupt bit
		CCHSNSB EQU X'04'	Sense data stored bit
		CCHCNTB EQU X'02'	Count valid bit
		CCHNRYB EQU X'01'	No retry bit
41	65	IGBLAME DS CL1	Probable source of error
		Bits defined in IGBLAME	
		CCHCPU EQU X'80'	CPU is source of error
		CCHCHNL EQU X'40'	Channel is source of error
		CCHSCUB EQU X'20'	Storage control unit is source of error
		CCHSTG EQU X'10'	Storage is source of error
		CCHINTFC EQU X'08'	I/O interface is source of error
42	66	IGVALIDB DS CL1	Validity indicator bits
		Bits defined in IGVALIDB	
		CCHRCV EQU X'10'	Retry code valid
		CCHUSV EQU X'08'	Selective reset
		CCHCMDV EQU X'04'	Command address valid
		CCHCAV EQU X'02'	Channel address valid
		CCHDAV EQU X'01'	Device address valid
43	67	IGTERMSQ DS CL1	Termination/sequence code bits
		Bits defined in IGTERMSQ	
		COMP SYS EQU X'C0'	System reset
		COMPSEL EQU X'80'	Selective reset
		COMPFES EQU X'40'	Forced ending sequence
		COMPID EQU X'00'	Interface disconnect
		CCHDI EQU X'08'	Disconnect in sequence code bits
		Sequence code bits	
		RTCODE0 EQU X'00'	Retry
		RTCODE1 EQU X'01'	Code
		RTCODE2 EQU X'02'	Values
		RTCODE3 EQU X'03'	For
		RTCODE4 EQU X'04'	The
		RTCODE5 EQU X'05'	Constructed
		RTCODE6 EQU X'06'	ECSW
		RTCODE7 EQU X'07'	
44	68	CCDEVTYP DS 1F	CP device type
48	72	CCHANID DS XL2	Channel ID
4A	74	CCHCUA DS 1H	Actual failing address
4C	76	CCHMP DS 1F	MP information
50	80	CCHLOG80 DS 0CL112	2880 channel - 112 bytes
50	80	CCHLOG70 DS 0CL24	2870 channel - 24 bytes
50	80	CCHLOG60 DS 0CL24	2860 channel - 24 bytes
		CCHSIZE1 EQU (*-CCHREC)/8	Size in doublewords (X'0A')



50	80	CUAADDR DS	CL4	Unit address stored by integrated channel
		CCHSIZE EQU	(*CCHREC)/8	Size in doublewords
54	84	CCHLOG45 DS	0CL96	145 integrated channel - 96 bytes
54	84	CCHLOG35 DS	0CL24	135 integrated channel - 24 bytes

CHXBLOK AND CHYBLOK - VIRTUAL CHANNEL-TO-CHANNEL ADAPTER CONTROL BLOCKS

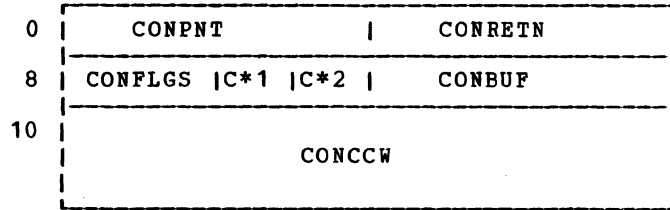
0	CHXOTHR		CHYOTHR
8	X*1  X*2  X*3  X*4  Y*1  Y*2  Y*3  Y*4		
10	CHXNCCW		CHYNCCW
18	CHXRCNT		CHYRCNT
20	CHXSTAT		CHYSTAT
	CHXYADD		CHXYADD
28	CHXIDAW		CHYIDAW
30	CHXCNCT		CHYCNCT
38	CHXWRK1		CHYWRK1

Displacement		Field	Field Description, Contents, Meaning	
Hex	Dec	Name		
		CHXBLOK		X-side channel adapter block
0	0	CHXOTHR DS 2F		VMBLOK address of Y-side adapter user
8	8	CHXFLAG DS 1X		X*1 Internal processing flags
		Bits defined in CHXFLAG and CHYFLAG:		
		CHBMNOP EQU X'80'		Modified NOP issued (also in CMDT)
		CHBM370 EQU X'40'		CTCA operating in System/370 mode
		CHBATTN EQU X'20'		Attention pending from Y-side
		CHBREST EQU X'10'		CTCA has been reset X-side and Y-side
		CHBEOFL EQU X'08'		Force EOF to next READ
9	9	CHXCMDB DS 1X		X*2 Active CCW command byte buffer
A	10	CHXCMDT DS 1X		X*3 Active CCW command type (RD, WR, etc.)
		Bits defined in CHXCMDT and CHYCMDT:		
		CHBCTNL EQU X'40'		Control, other than NOP
		CHBRDBK EQU X'20'		Read backward
		CHBWEOP EQU X'10'		Write EOF
		CHBSCMD EQU X'08'		Sense command byte
		CHBSADS EQU X'04'		Sense adapter status
		CHBREAD EQU X'02'		Read
		CHBWRT EQU X'01'		Write
B	11	CHXPKEY DS 1X		X*4 Virtual CAW protection key
C	12	DS 4X		
10	16	CHXNCCW DS 2F		Next CCW fetch address (real)
18	24	CHXRCNT DS 2F		Remaining CCW data count

20	32	CHXSTAT	DS	1H	Device status accumulation field
22	34	CHXYADD	DS	1H	Virtual address of Y-side adapter
24	36		DS	2H	
28	40	CHXIDAW	DS	2F	Active indirect-data-list word
30	48	CHXCNCT	DS	2F	CPEXBLOK for channel reconnect
38	56	CHXWRK1	DS	2F	Work area word
I		CHBSIZE	EQU	(*-CHXBLOK)/8	Total block size in doublewords (X'08')
		CHYBLOK	DSECT	,	Y-side channel adapter block
0	0	CHYOTHR	DS	2F	VMBLOK address of X-side adapter user
8	8	CHYFLAG	DS	1X	Y*1 Internal processing flags
Bits defined in CHXFLAG and CHYFLAG:					
		CHBMNOP	EQU	X'80'	Modified NOP issued (also in CMDT)
		CHBM370	EQU	X'40'	CTCA operating in System/370 mode
		CHBATTN	EQU	X'20'	Attention pending from Y-side
		CHBREST	EQU	X'10'	CTCA has been reset X-side and Y-side
		CHBEOFL	EQU	X'08'	Force EOF to next READ
9	9	CHYCMDDB	DS	1X	Y*2 Active CCW command byte buffer
A	10	CHYCMDT	DS	1X	Y*3 Active CCW command byte
Bits defined in CHXCMDT and CHYCMDT:					
		CHBCNTL	EQU	X'40'	Control, other than NOP
		CHBRDBK	EQU	X'20'	Read backward
		CHBWEOP	EQU	X'10'	Write EOF
		CHBSCMD	EQU	X'08'	Sense command byte
		CHBSADS	EQU	X'04'	Sense adapter status
		CHBREAD	EQU	X'02'	Read
		CHBWRT	EQU	X'01'	Write
B	11	CHYPKEY	DS	1X	Y*4 Virtual CAW protection key
C	12		DS	4X	
10	16	CHYNCCW	DS	2F	Next CCW fetch address
18	24	CHYRCNT	DS	2F	Remaining CCW data count
20	32	CHYSTAT	DS	1H	Device status accumulation field
22	34	CHYXADD	DS	1H	Virtual address of X-side adapter
24	36		DS	2H	
28	40	CHYIDAW	DS	2F	Active indirect-data-list word
30	48	CHYCNCCT	DS	2F	CPEXBLOK for channel reconnect
38	56	CHYWRK1	DS	2F	Work area word

**Note:** As indicated in the illustrated block, the CHXBLOK and CHYBLOK are interleaved with a 4-byte displacement. The X-side VDEVBLOK points to the +0 slot, the Y-side VDEVBLOK points to the +4 slot; however, once the virtual connection is made, either side can be the X-side or Y-side since this interleaved arrangement makes the control block references completely symmetrical. The dual DSECT definition allows the active adapter (defined to be the X-side, arbitrarily) to reference both adapter sides concurrently without knowing which is at +0 or +4.

CONTASK - CONSOLE I/O



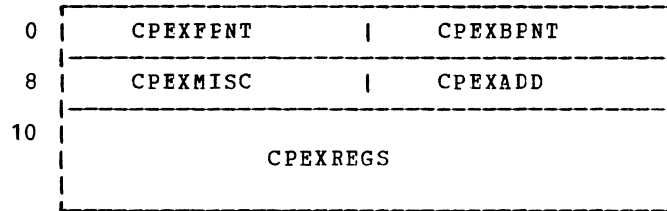
Displacement		Field		Field Description, Contents, Meaning
Hex	Dec	Name		
0	0	CONPNT	DS 1F	Pointer to next CONTASK
4	4	CONRETN	DS 1F	Pointer to CPEXBLOK for return
8	8	CONFLGS	DS 1H	CONTASK flags
A	10	CONTSKSZ	DS 1X	C*1 - CONTASK size in doublewords
B	11	CONBUFSZ	DS 1X	C*2 - BUFFER size in doublewords
C	12	CONBUF	DS 1F	Address of data BUFFER
10	16	CONCCW	DS 1D	One or more CCWs for console I/O
		CONTSIZE	EQU (CONCCW-CONTASK)/8	CONTASK size in doublewords (X'03')
			ORG CONFLGS	
8	8	CONSTAT	DS 4B	CONTASK Status
		Bits Defined in CONSTAT		
		CONOUTPT	EQU X'80'	Output CONTASK
		CONBUFVD	EQU X'40'	CONBUF contains a valid Free Storage Buffer
			ORG CONFLGS	
8	8	CONPARM	DS 1H	QUECONS parameter flags
			ORG CONCCW	
10	16	CONADDR	DS 1F	CCW data address
14	20	CONFLAG	DS 1X	CCW flag bits
15	21	CONRSV3	DS 1X	Reserved for IBM use
16	22	CONCNT	DS 1H	CCW byte count
			ORG CONADDR	
10	16	CONCOMND	DS 1X	CCW command code

CORTABLE - STORAGE ALLOCATION TABLE

0	CORFPNT		CORBPNT
8	C*1   CORSWPNT		CORPGPNT

Displacement Hex Dec	Field Name	Field Description, Contents, Meaning
0 0	CORFPNT DS 1F	Pointer to next CORETABLE entry in queue
4 4	CORBPNT DS 1F	Pointer to previous CORETABLE entry in queue
8 8	CORSWPNT DS 1F	Pointer to SWAPTABLE for page
C 12	CORPGPNT DS 1F	Pointer to PAGTABLE for page
8 8	CORFLAG DS 1X Bits Defined in CORFLAG CORIOLCK EQU X'80' CORCFLCK EQU X'40' CORFLUSH EQU X'20' CORFREE EQU X'10' CORSHARE EQU X'08' CORRSV EQU X'04' CORCP EQU X'02' CORDISA EQU X'01'	C*1 - CORTABLE entry status flags Page locked for I/O, CORLCNT greater than 0 Page locked by console function Page is in FLUSH list Page is in FREE list Page is shared Page is reserved Page belongs to CP Page disabled, not available
	Entry definition if page is locked	
4 4	CORLCNT DS 1F CORBPNT DS 1F	Page lock count for CORIOLCK
	Entry definition if page is in transit	
8 8	CORCODE DS 1X	C*1 - DASD op-code for PAGEIO

CPEXBLOK - CP EXECUTE BLCK



Displacement		Field			Field Description, Contents, Meaning
Hex	Dec	Name			
0	0	CPEXFPNT	DS	1F	Pointer to next CPEXBLOK
4	4	CPEXBPNT	DS	1F	Pointer to previous CPEXBLOK
8	8	CPEXMISC	DS	1F	Use varies with stacker
C	12	CPEXADD	DS	1F	Execute address
10	16	CPEXREGS	DS	16F	Execute registers

CPEXSIZE EQU (\*-CPEXBLOK)/8 Size in doublewords (X'0A')

For CPEXREGS Area  
ORG CPEXREGS

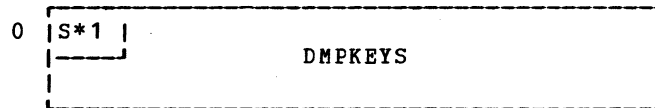
10	16	CPEXR0	DS	1F
14	20	CPEXR1	DS	1F
18	24	CPEXR2	DS	1F
1C	28	CPEXR3	DS	1F
20	32	CPEXR4	DS	1F
24	36	CPEXR5	DS	1F
28	40	CPEXR6	DS	1F
2C	44	CPEXR7	DS	1F
30	48	CPEXR8	DS	1F
34	52	CPEXR9	DS	1F
38	56	CPEXR10	DS	1F
3C	60	CPEXR11	DS	1F
40	64	CPEXR12	DS	1F
44	68	CPEXR13	DS	1F
48	72	CPEXR14	DS	1F
4C	76	CPEXR15	DS	1F

DMPINREC - DUMP FILE INFORMATION RECORD

0	DMPGPRS
40	DMPCRS
80	DMPFPRS
A0	DMPTODCK
A8	DMPCPUTM
B0	DMPCKCOM
B8	S*1   S*2   DMPRSV2   DMPSYSRV
C0	DMPLCORE
1C0	DMPPGMAP

Displacement		Field			Field Description, Contents, Meaning
Hex	Dec	Name			
0	0	DMPGPRS	DS	16F	16 General Registers
40	64	DMPCRS	DS	16F	16 Control Registers
80	128	DMPFPRS	DS	4D	4 Floating Point Registers (if Floating-Point Feature is installed.)
A0	160	DMPTODCK	DS	1D	Time-of-day clock
A8	168	DMPCPUTM	DS	1D	CPU timer
B0	176	DMPCKCOM	DS	1D	Time-of-day clock comparator
B8	184	DMPFLAG	DS	1X	S*1 - flag byte
		Bits Defined in DMPFLAG			
		HALFPAGE	EQU	X'80'	Last record in DUMP file = 2K
B9	185	DMPRSV1	DS	1X	S*2 - reserved for IBM use
BA	186	DMPRSV2	DS	1H	Reserved for IBM use
BC	188	DMPSYSRV	DS	1F	System generated storage size
C0	192	DMPLCORE	DS	256X	Locations 0-256 of storage memory
1C0	448	DMPPGMAP	DS	4096B	Bit map indicating which pages appear in the DUMP file (each bit represents a 4K block)

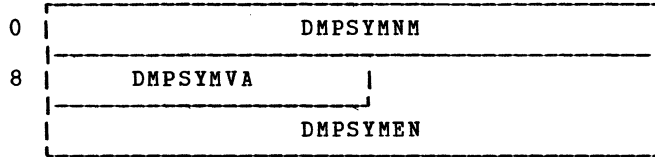
DMPKYREC - DUMP FILE KEY RECORD



Displacement		Field		Field Description, Contents, Meaning
Hex	Dec	Name		
0	0	DMPKEYS	DS 4096X	Main storage keys
0	0	DMPKEY	ORG DMPKEYS DS 1X	S*1 - storage key for each 2K block



DMPTBREC - DUMP FILE SYMEOI TABLE RECORD



Displacement		Field		Field Description, Contents, Meaning
Hex	Dec	Name		
0	0	DMP SYMEN DS	341XL12	Symbol table entries
			ORG DMP SYMEN	
0	0	DMP SYMNM DS	CL8	CSECT or entry point name
8	8	DMP SYMVA DS	A	Location in main storage of this symbol

ECBLCK - EXTENSION TO VMPLCK FOR VIRTUAL MACHINE WITH RELOCATE WITH RELOCATE

0	EXTCR0		EXTCR1		
8	EXTCR2		EXTCR3		
10	EXTCR4		EXTCR5		
18	EXTCR6		EXTCR7		
20	EXTCR8		EXTCR9		
28	EXTCR10		EXTCR11		
30	EXTCR12		EXTCR13		
38	EXTCR14		EXTCR15		
40	EXTSHCR0		EXTSHCR1		
48	EXTSHLEN		EXTCPLEN		EXTCOPY
50	EXTSHSEG		EXTSEGLN		EXTARCH
58	EXTPERAD		EXTPERCD		EXTRSV1
60	EXTCPTMR				
68	EXTCPTRQ		EXTCCTRQ		

Displacement		Field		Field Description, Contents, Meaning
Hex	Dec	Name		
0	0	EXTCR0	DS 1F	Virtual control register 0; architecture controls Virtual control register 1; segment table pointer Virtual control register 2 ...thru register 15
4	4	EXTCR1	DS 1F	
8	8	EXTCR2	DS 1F	
C	12	EXTCR3	DS 1F	
10	16	EXTCR4	DS 1F	
14	20	EXTCR5	DS 1F	
18	24	EXTCR6	DS 1F	
1C	28	EXTCR7	DS 1F	
20	32	EXTCR8	DS 1F	
24	36	EXTCR9	DS 1F	
28	40	EXTCR10	DS 1F	
2C	44	EXTCR11	DS 1F	
30	48	EXTCR12	DS 1F	

34	52	EXTCR13	DS	1F	
38	56	EXTCR14	DS	1F	
3C	60	EXTCR15	DS	1F	
40	64	EXTSHCR0	DS	1F	Shadow control register 0
44	68	EXTSHCR1	DS	1F	Shadow control register 1
48	72	EXTSHLEN	DS	1H	Length of shadow SEGTABLE in bytes
4A	74	EXTCPLEN	DS	1H	Length of copy SEGTABLE in bytes
4C	76	EXTCOPY	DS	1F	Pointer to copy segment table
50	80	EXTSHSEG	DS	1F	Real address of shadow SEGTABLE
54	84	EXTSEGLN	DS	1H	Length of shadow SEGTABLE in doublewords
56	86	EXTARCH	DS	1H	Architecture control index
58	88	EXTPERAD	DS	1F	PER interrupt address
5C	92	EXTPERCD	DS	1H	PER interrupt code to be reflected
5E	94	EXTRSV1	DS	1H	Reserved for IBM use
60	96	EXTCPTMR	DS	1D	Virtual CPU timer
68	104	EXTCPTRQ	DS	1F	Address of TRQBLOK for CPU timer
6C	108	EXTCCTRQ	DS	1F	Address of TRQBLOK for clock comparator
		EXTSIZE	EQU	(*-ECBLOK)/8	ECBLOK size in doublewords (X'0E')

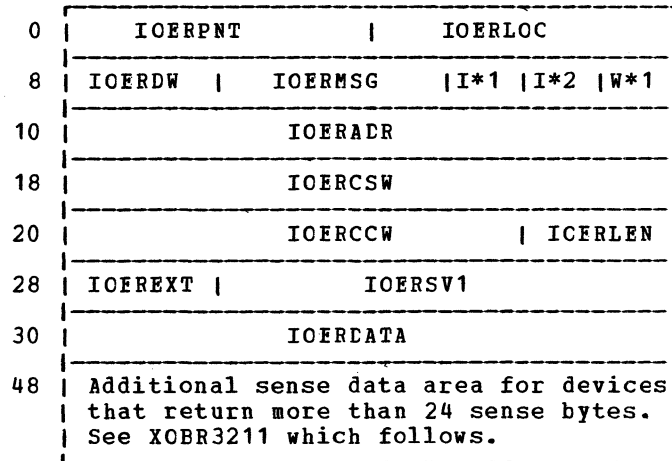
IOBLOK - I/O CONTROL BLOCK

0	IOBRADD	I*1	I*2		IOBLINK	
8	IOBFPNT				IOBEPNT	
10	IOBCYL	IOBVADE			IOBMISC	
18	IOBUSER				IOBIRA	
20	IOBCAW				IOBRCAW	
28		IOECSW				
30	IOBIOFR				IOBMISC2	
38	I*3	V*1		V*2		V*3

Displacement		Field			Field Description, Contents, Meaning
Hex	Dec	Name			
0	0	IOBRADD	DS	1H	Real device address for SIO
2	2	IOBFLAG	DS	1X	I*1 - IOBLOK flags
		Bits Defined in IOBFLAG			
		IOBCP	EQU	X'80'	CP generated I/O operation
		IOBRSTRT	EQU	X'40'	Restarted operation - IOBRCAW
		IOBSPLI	EQU	X'20'	DASD - CP split seek operation
		IOBPAG	EQU	X'10'	IOBLOK created for paging I/O
		IOBRELCU	EQU	X'08'	Control unit released at initiation
		IOBERP	EQU	X'04'	I/O task is under control of ERP
		IOBRES	EQU	X'02'	I/O task has been reset
		IOBHVC	EQU	X'01'	I/O initiated via DIAGNOSE instruction
3	3	IOBSTAT	DS	1X	I*2 - IOBLOK status
		Bits Defined in IOBSTAT			
		IOBFATAL	EQU	X'80'	Uncorrectable error in this I/O operation
		IOBUC	EQU	X'40'	Unit check status
		IOBSNSIO	EQU	X'20'	Sense operation (IOBSNSE)
		IOBREQUE	EQU	X'10'	Restarted operation (IOBCAW)
		IOBWRAF	EQU	X'08'	I/O task for autopoll wrap list
		IOBCC0	EQU	X'00'	Processing I/O interrupt
		IOBCC1	EQU	X'01'	Processing CC 1, CSW stored
		IOBCC2	EQU	X'02'	Processing CC 2, channel busy
		IOBCC3	EQU	X'03'	Processing CC 3, not available
4	4	IOBLINK	DS	1F	Reserved for IBM use

8	8	IOBFPNT	DS	1F	Pointer to next IOBLOK in queue
C	12	IOBBPNT	DS	1F	Pointer to previous IOELOK in queue
		IOBMSIZE	EQU	(*IOBLOK)/8	Multiple path IOBLOK size in dbl. wds (X'02')
10	16	IOBCYL	DS	1H	DASD - seek cylinder for this IOBLOK
12	18	IOBVADD	DS	1H	Virtual device address
14	20	IOBMISC	DS	1F	Use varies according to caller
18	24	IOBUSER	DS	1F	Pointer to VMBLOK of user
1C	28	IOBIRA	DS	1F	IOBLOK interrupt return address
20	32	IOBCAW	DS	1F	Pointer to CCW chain
24	36	IOBRCAW	DS	1F	Pointer to restart CCW chain
28	40	IOBCSW	DS	1D	Real CSW for I/O operation
30	48	IOBIOER	DS	1F	Pointer to IOERBLOK with sense
34	52	IOBMISC2	DS	1F	Use varies according to caller
38	56	ICBSPEC	DS	1X	I*3 - IOBLOK special requests
		Bits Defined in IOBSPEC			
		IOBTIO	EQU	X'80'	IOBLOK request for a 'TIO'
		IOBHIO	EQU	X'40'	IOBLOK request for a 'HIO'
39	57	IOBSV1	DS	1X	V*1 reserved for IBM use
3A	58	IOBSV2	DS	XL2	V*2 reserved for IBM use
3C	60	IOBSV3	DS	1F	V*3 reserved for IBM use
		IOBSIZE	EQU	(*IOBLOK)/8	IOBLOK size in doublewords (X'08')
		For CP IOBLCKs			
			ORG	IOBVADD	
12	18	IOBRCNT	DS	1H	Retry count

IOERBLOK - I/C ERRCR INFORMATION BLOK



Displacement		Field		Field Description, Contents, Meaning
Hex	Dec	Name		
0	0	IOERPNT	DS 1F	Pointer to next IOERBLOK
4	4	IOERLOC	DS 1F	Address of CCWs used in recovery
8	8	IOERDW	DS 1H	Size in doublewords of storage to construct CCWs
A	10	IOERMSG	DS XL3	Communications with ERP and message writer
			ORG IOERMSG	
A	10	IOERNUM	DS 1X	Message number for message writer
B	11	IOERIND3	DS 1X	Indicators for message writer
		Bits Defined in IOERIND3		
		IOERIGN	EQU X'80'	Allow IGNORE response
		IOERETRY	EQU X'40'	Allow RETRY response
		IOERCAN	EQU X'20'	Allow CANCEL response
		IOEREC	EQU X'10'	Error occurred during recovery action
		IOERDASD	EQU X'08'	Home address is present
		IOERDEC	EQU X'04'	Operator decision is necessary
		ICERINFO	EQU X'02'	Informational message
		IOERACT	EQU X'01'	Operator action is required
C	12	IOERIND4	DS 1X	Indicators for message writer
		Bits Defined in IOERIND4		
		IOERIGNR	EQU X'80'	Operator responded IGNORE
		IOERSTRT	EQU X'40'	Operator responded RETRY

D 13

IOERFLG1 DS 1X  
 Bits Defined in IOERFLG1  
 IOERPENDEQU X'80'  
 IOERCLNEQU X'40'  
 ICERERPEQU X'40'  
 IOERFSREQU X'20'  
 IOERDEPDEQU X'20'  
 IOERBSREQU X'10'  
 IOERDERDEQU X'10'  
 IOERERGEQU X'08'  
 IOERORAEQU X'04'  
 IOERSUPPEQU X'02'  
 IOERVLD EQU X'01'

I\*1 - IOERFLG1 field  
 Pending device end interrupt from interrupt request  
 Tape cleaning in progress  
 Spooling - error routine in control  
 Forward space record being executed  
 Spooling - waiting for device end  
 Backspace record being executed  
 Spooling - device end received  
 Erase gap command in progress  
 Opposite recovery action in progress  
 CCW has suppress data transfer bit on  
 Read opposite recovery successful

E 14

IOERFLG2 DS 1X  
 Bits Defined in IOERFLG2  
 IOERSTAT EQU X'80'  
 IOERHAEQU X'40'  
 IOERCAL EQU X'20'  
 IOERECF EQU X'10'  
 IOERREW EQU X'08'  
 IOERCYLR EQU X'04'  
 IOERCEMD EQU X'02'

I\*2 - IOERFLG2 field  
 Statistical data being unloaded  
 DASD home address being read  
 Stand alone recalibrate being executed  
 Error correction function  
 Tape rewind being executed  
 Cylinder (in sense byte) has been relocated  
 Intensive recording mode

F 15  
 10 16  
 18 24  
 20 32

IOERWRK DS 1X  
 IOERADR DS 1D  
 ICERCSW DS 1D  
 IOERCCW DS 1D

W\*1 - Miscellaneous work area  
 Home address for DASD devices  
 CSW associated with error  
 Sense CCW used to sense the real device

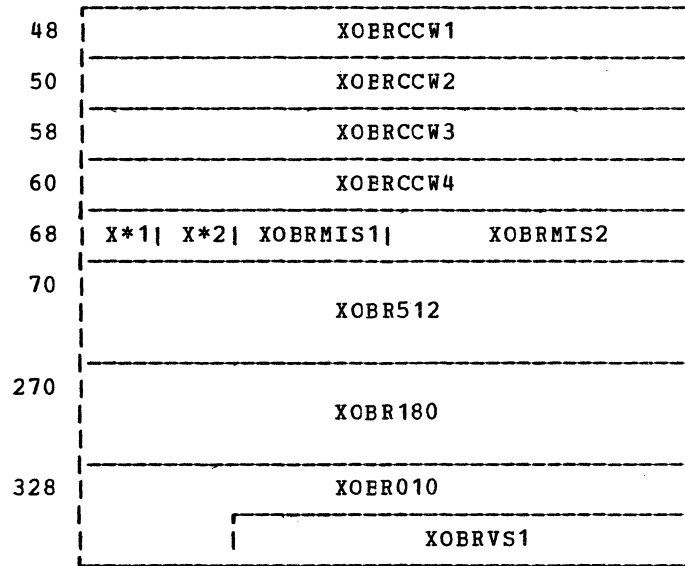
26 38  
 28 40  
 2A 42  
 30 48

ORG IOERCCW+6  
 IOERLEND S 1H  
 IOEREXT DS 1H  
 IOERSV1 DS XL6  
 IOERDATA DS 3D

Number of sense bytes present  
 Size of extended sense area in doublewords  
 Reserved for IBM use  
 Sense bytes associated with error

IOERSIZE EQU (\*-IOERBLOK)/8 IOERBLOK size in doublewords (X'09')

IOERBLOK DSECT CCNTINUE  
 X0BR3211 - EXTENDED OUTBOARD RECORDING BLOCK

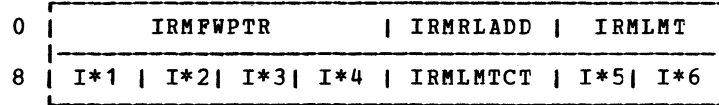


Displacement Hex Dec	Field Name		Field Description, Contents, Meaning
48 72	XOBRCCW1 DS	1D	CCW used to read OBR information
50 80	XOBRCCW2 DS	1D	CCW used to read OBR information
58 88	XOBRCCW3 DS	1D	CCW used to read OBR information
60 96	XOBRCCW4 DS	1D	CCW used to read OBR information
68 104	XOBRFLAG DS	1X	X*1 - XOBRFLAG field
	Bits Defined in XOBRFLAG		
	XOBRRT1 EQU	X'80'	T1 Buffer type information present
	XOBRRT2 EQU	X'40'	T2 Buffer type information present
	XOBRRT3 EQU	X'20'	T3 Buffer type information present
69 105	XOBRSTAT DS	1X	X*2 - XOBRSTAT field
	Bits defined in XOBRSTAT		
	XOBRRT1 EQU	X'80'	Perform routine 1 in error module
	XOBRRT2 EQU	X'40'	Perform routine 2 in error module
	XOBRRT3 EQU	X'20'	Perform routine 3 in error module
	XOBRRT4 EQU	X'10'	Perform routine 4 in error module
	XOBRRT5 EQU	X'08'	Perform routine 5 in error module



		XOBRRT6	EQU	X'04'	Perform routine 6 in error module	
		XOBRRT7	EQU	X'02'	Perform routine 7 in error module	
		XOBRRT8	EQU	X'01'	Perform routine 8 in error module	
	6A	106	XOBRMIS1	DS	1H	Used by the error routine
	6C	108	XOBRMIS2	DS	1F	Used by the error routine
	70	112	XOBR512	DS	CL512	Space for USCB data
	270	624	XOBR180	DS	CL184	Space for FCB data
				ORG	XOBR180	
	270	624	XOBR150	DS	CL150	Space for PLB check data
				ORG		
	328	808	XOBR010	DS	CL10	Space for first ten error characters
	332	818	XOBRSV1	DS	CL6	Reserved for IBM use
			XOBRSIZE	EQU	(*-IOERBLOK)/8	Size of IOER and XOBR in double words (X'67')
			XOBREXT	EQU	(*-XOBRCCW1)/8	Size of XOBR3211 in double words (X'5E')

IRMBLOCK - INTENSIVE ERROR RECORDING MODE BLOCK



Displacement		Field		Field Description, Contents, Meaning
Hex	Dec	Name		
0	0	IRMFWPTR DS	1F	Reserved for IBM use
4	4	IRMLADD DS	1H	Device address
6	6	IRMLMT DS	1H	Limit count - every 'nth' record is requested.
8	8	IRMBYT1 DS	1X	I*1 - first sense byte specified
9	9	IRMBIT1 DS	1X	I*2 - sense bit within first sense byte
A	10	IRMEYT2 DS	1X	I*3 - second sense byte specified
B	11	IRMBIT2 DS	1X	I*4 - sense bit within second sense byte
C	12	IRMLMTCT DS	1H	Temporary summary count for limit detection
E	13	IRMMAXCT DS	1X	I*5 - count of recordings made for this request
F	15	IRMFLG DS	1X	I*6 - flag byte
Bits Defined in IRMFLG				
		IRMAND EQU	X'80'	AND condition specified
		IRMOR EQU	X'40'	OR condition specified
		IRMSIZE EQU	(*-IRMBLOCK)/8	IRMBLOCK size in doublewords (X'02')

MCHAREA - MACHINE CHECK SAVE AREA

0	MCDAMLEN	MCHRESEV
8	M*1	M*2   M*3   M*4   M*5   M*6   M*7   M*8
10	MCHLSUM	
38	N*1	N*2   N*3   N*4   N*5   N*6   N*7   N*8
40	MCHFSAR	MCHFSAV
48	MCHFSEAV	MCHPDARI
50	L*1	L*2   L*3   L*4   CPULIMIT   MCHRES1
58	BUFDIA55	BUF55DIA
60	BUFENA55	BUF55ENA
68	ECCDIS55	ECC55DIS
70	ECCENA55	ECC55ENA
78	BUFDIA65	
80	BUFENA65	
88	ECCDIS65	
90	ECCENA65	

Displacement  
Hex Dec

Field  
Name

Field Description, Contents, Meaning

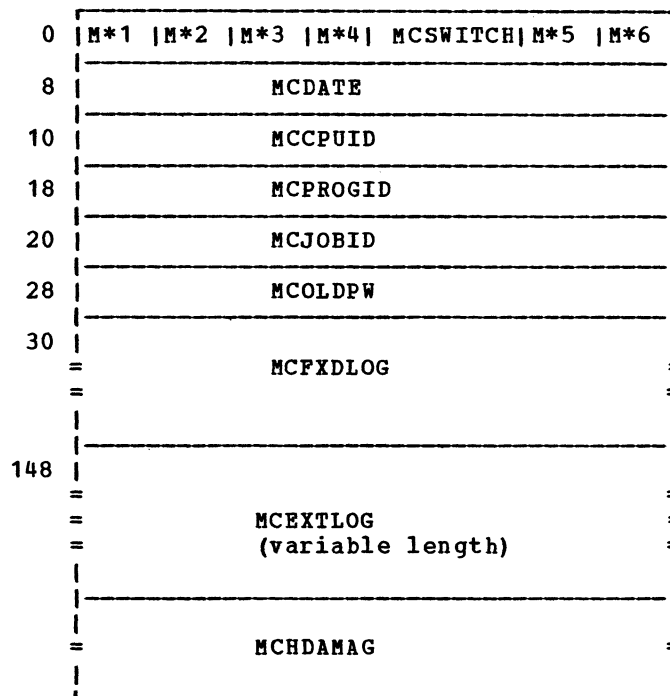
Hex	Dec	Field Name	Field Description, Contents, Meaning
0	0	MCDAMASS DS 0D	Damage assessment
0	0	MCDAMLEN DS 1H	Length of the damage assessment field
2	2	MCHRESEV DS XL6	Reserved for IBM use
8	8	MCHDAMFL DS 0BL8	Damage assessment data
8	8	MCHFLAGO DS 1X	M*1 System status
		Bits defined in MCHFLAGO	
		MCHOHDWR EQU X'80'	Hardware recovery
		MCHOSFTR EQU X'40'	Software recovery
		MCHOUSAD EQU X'20'	User aborted

		MCH0TERM EQU	X'08'	Operating system termination
		MCH0QUIT EQU	X'04'	Quiet mode in effect
9	9	MCHFLAG1 DS	1X	M*2 Damage area
		Bits defined in MCHFLAG1		
		MCH1MAIN EQU	X'80'	Main storage
		MCH1BUFF EQU	X'40'	Buffer
		MCH1COST EQU	X'20'	Control storage
		MCH1PROC EQU	X'08'	Processor
		MCH1TODC EQU	X'02'	Time-of-day clock
		MCH1SYS D EQU	X'01'	System damage
A	10	MCHFLAG2 DS	1X	M*3 Damage area (continued)
B	11	MCHFLAG3 DS	1X	M*4 Error type
		Bits defined in MCHFLAG3		
		MCH3INTE EQU	X'80'	Intermittent
		MCH3SOLD EQU	X'40'	Solid
		MCH3DATA EQU	X'20'	Data
		MCH3PROT EQU	X'10'	Protect
C	12	MCHFLAG4 DS	1X	M*5 RMS Action data
		Bits defined in MCHFLAG4		
		MCH4TOLO EQU	X'80'	Time out loop
		MCH4REPA EQU	X'40'	Repair
		MCH4STRE EQU	X'20'	Storage reconfigure
		MCH4BURE EQU	X'10'	Buffer reconfigure
D	13	MCHFLAG5 DS	1X	M*6 RMS Information status
		Bits defined in MCHFLAG5		
		MCH5INLG EQU	X'80'	Invalid logout
		MCH5INMC EQU	X'40'	Invalid machine check interrupt code
		MCH5IFSA EQU	X'20'	Invalid failing storage address
E	14	MCHFLAG6 DS	1X	M*7 RMS wait state suffix
F	15	MCHFLAG7 DS	1X	M*8 RMS information status
		Bits defined in MCHFLAG7		
		MCH7SMCR EQU	X'80'	Second machine check recursion
		MCH7VRTM EQU	X'40'	Terminate the virtual user
		MCH7OPSW EQU	X'10'	M.C. old PSW in problem state
		MCH7VEQR EQU	X'08'	Terminate the Virtual equal Real user
10	16	MCHLSUM DS	1X	Summary
38	56	MCHPDAR DS	OBL8	
38	56	MCHPDAR0 DS	1X	N*1 Action taken
39	57	MCHPDAR1 DS	1X	Failure type
		Bits defined in MCHPDAR1		
		MCHP1SDE EQU	X'80'	Solid storage data error
		MCHP1IDE EQU	X'40'	Intermittent storage data error
		MCHP1SKE EQU	X'20'	Solid SPF key error
		MCHP1IKE EQU	X'10'	Intermittent SPF key error

3A	58	MCHPDAR2 DS	1X	N*3 Operating system status
3B	59	MCHPDAR3 DS	1X	N*4 Location of failure
3C	60	MCHPDAR4 DS	1X	N*5 Location of failure
3D	61	MCHPDAR5 DS	1X	N*6 Requested operator awareness
3E	62	MCHPDAR6 DS	1X	Footprint
		Bits defined in MCHPDAR6		
		MCHP6CBA EQU	X'80'	Change bit active
3F	63	MCHPDAR7 DS	1X	Footprints
		Bits defined in MCHPDAR7		
		MCH7STCK EQU	X'80'	Interfaces for STACK routine
		MCH7GSTR EQU	X'40'	Interfaces for GETMAIN routine
		MCH7PURG EQU	X'20'	Interfaces for PURGE routine
		MCH7LOGO EQU	X'10'	Interfaces for V=R LOGOFF routine
		MCH7EXIT EQU	X'08'	Interfaces for exit to CP
		MCH7RSRE EQU	X'04'	Interfaces for RELEASE and RESET routines
		MCH7IOEM EQU	X'02'	Interfaces for the recorder
40	64	MCHFSAR DS	1F	Failing location real address
44	68	MCHPSAV DS	1F	Instruction address at failure
48	72	MCHFSEAV DS	1F	End of the failing location
4C	76	MCHPDARI DS	1F	End of failing storage address - virtual
		MCHLEN1 EQU	*-MCDAMASS	Length of damage assessment area
		MCHLEN EQU	*-MCHRESEV	Length of area to be cleared
50	80	MCHMODEL DS	1X	L*1 The model number for the machine
		Bits defined in MCHMODEL		
		NOMODEL EQU	X'00'	No support for machine
		MODEL135 EQU	X'04'	ID number for the 135 machine
		MODEL145 EQU	X'08'	ID number for the 145 machine
		MODEL155 EQU	X'0C'	ID number for the 155 machine
		MODEL158 EQU	X'0C'	ID number for the 158 machine
		MODEL165 EQU	X'10'	ID number for the 165 machine
		MODEL168 EQU	X'10'	ID number for the 168 machine
51	81	SWITCH DS	1X	L*2 Main storage exercise switch
52	82	MODEFLAG DS	1X	L*3 Flag field for MODE command
		Bits defined in MODEFLAG		
		MODEQUIT EQU	X'80'	ECC is in QUIET mode
53	83	MODFLAG1 DS	1X	L*4 Flag field for message indicator in MODE command
		Bits defined in MODFLAG1		
		MOD1RETY EQU	X'80'	Message indicator for RETRY message
		MOD1QUIT EQU	X'40'	Message indicator for QUIET message
54	84	CPULIMIT DS	1H	The count field for soft error
56	86	MCHRES1 DS	1H	Reserved for IBM use
58	88	DS	0D	

58	88	BUFDIA55	DC	X'0100D100'	Disable buffer for Model 155
5C	92	BUF55DIA	DS	1F	Reserved for IBM use
60	96	BUFENA55	DC	X'0200D100'	Enable buffer for Model 155
64	100	BUF55ENA	DS	1F	Reserved for IBM use
68	104	ECCDIS55	DC	X'0300D100'	Disable ECC for Model 155
6C	108	ECC55DIS	DS	1F	Reserved for IBM use
70	112	ECCENA55	DC	X'0400D100'	Enable ECC for Model 155
74	116	ECC55ENA	DS	1F	Reserved for IBM use
78	120	BUFDIA65	DC	X'0300000000000000'	Disable buffer for Model 165
80	128	BUFENA65	DC	X'0300002000000000'	Enable buffer for Model 165
88	136	ECCDIS65	DC	X'0200000003000000'	Disable ECC for Model 165
90	144	ECCENA65	DC	X'0200000000000000'	Enable ECC for Model 165
		MCHFIX	EQU	280+48	The length of the fixed logout and header record for machine check handler
		MCHLEN2	EQU	*-MCDAMASS	The communication area length

MCRECORD - MACHINE CHECK HANDLER RECORD



Displacement		Field		Field Description, Contents, Meaning
Hex	Dec	Name		
0	0	MCREC	DS 0D	
0	0	MCRECTYP	DS 1X	M*1 Machine check record type
1	1	MCOPSYS	DS 1X	M*2 Operating system
2	2	MCSWONE	DS 1X	M*3 Record independent switch
3	3	MCSWTWO	DS 1X	M*4 Record dependent switch
4	4	MCSWITCH	DS 2X	Unused switches
6	6	MCRECCNT	DS 1X	M*5 Record count
7	7	MCRECCC	DS 1X	M*6 Spare
8	8	MCDATE	DS XL8	Date and time
10	16	MCCPUID	DS XL8	CPU identification
18	24	MCPROGID	DS XL8	Program identity
20	32	MCJOBID	DS XL8	Job identity (unused)
28	40	MCOLDPW	DS XL8	Machine check old PSW
30	48	MCFXDLOG	DS 35D	Machine check fixed logout

	FXDLGLH EQU	(*MCFXDLOG)	
148 328	MCEXTLOG EQU	*	Machine check extended logout (the extended logout length is variable length — machine dependent)
	MCHDAMAG EQU	*	The damage assessment area (80 bytes)



OWNDLIST - CP OWNED VOLUMES LIST

0	OWNDVSER	OWNRDEV
---	----------	---------

Displacement		Field		Field Description, Contents, Meaning
Hex	Dec	Name		
0	0	OWNDVSER DS	CL6	Volume serial number
6	6	OWNRDEV DS	1H	Displacement of RDEVBLK for the volume
			ORG	
6	6	OWNDPREF DS	OWNRDEV 1X	Allocation preference

PAGTABLE - PAGE TABLE

0	PAGRSV1		PAGSWP
8	PAGCORE		

Displacement		Field		Field Description, Contents, Meaning
Hex	Dec	Name		
0	0	PAGRSV1 DS	1F	Reserved to align PAGCORE on a doubleword
4	4	PAGSWP DS	1F	Pointer to SWPTABLE
8	8	PAGCORE DS	1H	Real page address
		Bits Defined in PAGCORE+1		
		PAGINVAL EQU	X'08'	PAGTABLE entry invalid
		PAGREF EQU	X'01'	Page has been referenced

PSA - PREFIX STORAGE AREA (LOW STORAGE LOCATIONS)

Page 0, Machine Usage

0	IPLPSW		IPLCCW1
10	IPLCCW2		EXOPSW
20	SVCOPSW		PROPSW
30	MCOPSW		IOOPSW
40	CSW		CAW  QUANTUMR
50	TIMER  QUANTUM		EXNPSW
60	SVCNPSW		PRNPSW
70	MCNPSW		IONPSW
80	CPULOG		
100	FXDLOG		
160	FPRLOG		
180	GRLOG		
1C0	CRLOG		
200	TEMPSAVE		
240	BALRSAVE		
280	FREESAVE		
2C0	FREWORK		
2F0	DATE		TODATE

300	STARTIME		CPUID
310	IDLEWAIT		PAGEWAIT
320	IONTWAIT		PROBTIME
330	RUNPSW		RUNUSER  DSP LPSW
340	RUNCRO		RUNCR1  CPSTAT  CPRESTRT
350	PGREAD		PGWRITE   PGWAITIM
360	PGWAITPG		PSASVCCT  P*1  P*2
370	CPID		CPABEND  P*3  P*4  ASYVM
380	ARSPPR		ARSPPU  ARSPRD  ARIOPU
390	ARIOPR		ARIORD  PSARSV6  ARSPAC
3A0	AVMREAL		ASYSABND  ASYSLC  ASYSOP
3B0	ARIOCT		ARIOCH  ARIOCU  ARIODV
3C0	ARIOCC		ARIOUC  ARIODC  ACORETBL
3D0	APAGCP		CPCREG0  CPCREG8  PSARSV9
3E0	PSARSV10		PSARSV11  ADMKFVR  XVRINST
3F0	PAGECUR		MONNEXT  PAGEEND  PAGENXT
400	TRACEFLG		PSARSV12
PSARSV15			
430	INSTWRD1		INSTWRD2  INSTWRD3  INSTWRD4

Displacement		Field		Field Description, Contents, Meaning
Hex	Dec	Name		
0	0	IPLPSW	DS 1D	IPL start PSW
8	8	IPLCCW1	DS 1D	IPL CCW
10	16	IPLCCW2	DS 1D	IPL CCW
			ORG IPLCCW1	
8	8	PSARSV3	DS 1F	Reserved for IBM use
C	12	TRACSTRT	DS 1F	Pointer to start of trace table
10	16	TRACEND	DS 1F	Pointer to end of trace table
14	20	TRACCURR	DS 1F	Pointer to next available trace table entry
18	24	EXOPSW	DS 1D	External old PSW
20	32	SVCOPSW	DS 1D	SVC old PSW
28	40	PROPSW	DS 1D	Program old PSW
30	48	MCOPSW	DS 1D	Machine check old PSW
38	56	IOOPSW	DS 1D	I/O old PSW
40	64	CSW	DS 1D	Channel status word
48	72	CAW	DS 1F	Channel address word
4C	76	QUANTUMR	DS 1F	Interval timer value at last interrupt
50	80	TIMER	DS 1F	13 microsecond interval timer
54	84	QUANTUM	DS 1F	Interval timer value at last dispatch
58	88	EXNPSW	DS 1D	External new PSW
60	96	SVCNPSW	DS 1D	SVC new PSW
68	104	PRNPSW	DS 1D	Program new PSW
70	112	MCNPSW	DS 1D	Machine check new PSW
78	120	IONPSW	DS 1D	I/O new PSW
80	128	CPULCG	DS 16D	CPU and storage logout area
			ORG CPULOG	
80	128		DS 1F	Reserved for IBM use
84	132	INTEXF	DS 1F	External interrupt code (fullword)
86	134	INTEX	EQU INTEXF+2	External interrupt code (halfword)
88	136	INTSVCL	DS 1H	SVC instruction length code (ILC)
8A	138	INTSVC	DS 1H	SVC interrupt code
8C	140	INTPRL	DS 1H	Program instruction length code (ILC)
8E	142	INTPR	DS 1H	Program interrupt code
90	144	TREXADD	DS 1F	Translation exception address
94	148	MONCLASS	DS 1H	Monitor class
96	150	PERCODE	DS 1H	PER interrupt code
98	152	PERADD	DS 1F	PER interrupt address
9C	156	MONCODE	DS 1F	Monitor code
A0	160		DS 1D	Reserved for IBM use
A8	168	CHANID	DS 1F	Channel identification
AC	172	IOELPNTR	DS 1F	I/O extended logout (IOEL) pointer
E0	176	ECSWLOG	DS 1F	Limited channel logout (ECSW)
B4	180		DS 1F	Reserved for IBM use
E8	184	INTKFLIN	DS 1F	I/O interrupt key, flags, interface address
BA	186	INTTIO	EQU INTKFLIN+2	I/O interrupt device address (halfword)
EC	188		DS 11F	Reserved for IBM use

E8	232	INTMC	DS	1D	Machine check interrupt code
F0	240		DS	1D	Reserved for IBM use
F8	248	FAILSTAD	DS	1F	Failing storage address
FC	252	REGNCODE	DS	1F	Region code
100	256	FXDLOG	DS	12D	Fixed logout area
160	352	FPRLOG	DS	4D	Floating-point register logout area
180	384	GRLCG	DS	16F	General register logout area
1C0	448	CRLOG	DS	16F	Control register logout area
200	512	CPUSAGE	DS	0H	End of machine usage, start of CP usage
			ORG	CPUSAGE	
200	512	TEMPSAVE	DS	16F	Temporary save area
			ORG	TEMPSAVE	
200	512	TEMPRO	DS	1F	
204	516	TEMPR1	DS	1F	
208	520	TEMPR2	DS	1F	
20C	524	TEMPR3	DS	1F	
210	528	TEMPR4	DS	1F	
214	532	TEMPR5	DS	1F	
218	536	TEMPR6	DS	1F	
21C	540	TEMPR7	DS	1F	
220	544	TEMPR8	DS	1F	
224	548	TEMPR9	DS	1F	
228	552	TEMPR10	DS	1F	
22C	556	TEMPR11	DS	1F	
230	560	TEMPR12	DS	1F	
234	564	TEMPR13	DS	1F	
238	568	TEMPR14	DS	1F	
23C	572	TEMPR15	DS	1F	
240	576	BALRSAVE	DS	16F	BALR linkage save area
			ORG	BALRSAVE	
240	576	BALR0	DS	1F	
244	580	BALR1	DS	1F	
248	584	BALR2	DS	1F	
24C	588	BALR3	DS	1F	
250	592	BALR4	DS	1F	
254	596	BALR5	DS	1F	
258	600	BALR6	DS	1F	
25C	604	BALR7	DS	1F	
260	608	BALR8	DS	1F	
264	612	BALR9	DS	1F	
268	616	BALR10	DS	1F	
26C	620	BALR11	DS	1F	
270	624	BALR12	DS	1F	
274	628	BALR13	DS	1F	
278	632	BALR14	DS	1F	
27C	636	BALR15	DS	1F	
280	640	FREESAVE	DS	16F	DMKFRE save area

280	640	FREERO	DS	1F	
284	644	FREER1	DS	1F	
288	648	FREER2	DS	1F	
28C	652	FREER3	DS	1F	
290	656	FREER4	DS	1F	
294	660	FREER5	DS	1F	
298	664	FREER6	DS	1F	
29C	668	FREER7	DS	1F	
2A0	672	FREER8	DS	1F	
2A4	676	FREER9	DS	1F	
2A8	680	FREER10	DS	1F	
2AC	684	FREER11	DS	1F	
2B0	688	FREER12	DS	1F	
2B4	692	FREER13	DS	1F	
2B8	696	FREER14	DS	1F	
2BC	700	FREER15	DS	1F	
2C0	704	FREWORK	DS	12F	DMKPRE work area
2F0	752	DATE	DS	CL8	Date - mm/dd/yy - edited EBCDIC
2F8	760	TODATE	DS	1D	TOD clock at 00.00.00 today - local time
300	768	STARTIME	DS	1D	Date and time started - TOD clock value
308	776	CPUID	DS	1D	CPU identification
			ORG	CPUID	
308	776	CPUVERS	DS	1X	Version code
309	777	CPUSER	DS	3X	CPU serial number - packed unsigned
30C	780	CPUMODEL	DS	2X	CPU model number
30E	782	CPUMCELL	DS	1H	MAXIMUM length in bytes of MCEL
310	784	IDLEWAIT	DC	X'7FFFFFFFFFFFFFFF000'	Total system idle wait time
318	792	PAGWAIT	DC	X'7FFFFFFFFFFFFFFF000'	Total system page wait time
320	800	IONTWAIT	DC	X'7FFFFFFFFFFFFFFF000'	Total system I/O wait time
328	808	PROBTIME	DC	X'7FFFFFFFFFFFFFFF000'	Total system problem state time
330	816	RUNPSW	DS	1D	PSW last loaded by Dispatcher
338	824	RUNUSER	DS	1F	Address of dispatched VMBLOK
33C	828	DSPLPSW	DS	1F	Load PSW instruction used to dispatch
340	832	RUNCRO	DS	1F	Control register zero at dispatch
344	836	RUNCR1	DS	1F	Control register one at dispatch
348	840	CPSTAT	DS	1F	CP running status
			ORG	CPSTAT	
348	840	CPSTATUS	DS	1X	CP running status
		Bits defined in CPSTATUS			
		CPWAIT	EQU	X'80'	CP in wait state
		CPRUN	EQU	X'40'	CP running user in RUNUSER
		CPEX	EQU	X'20'	CP executing stacked request
		CPFVRUN	EQU	X'10'	Reserved for IBM use
34C	844	CPRESTR	DS	1F	Restart address if external interrupt marks page invalid
350	848	PGREAD	DS	1F	Total number of page reads
354	852	PGWRITE	DS	1F	Total number of page writes

358	856	PGWAITIM	DS	1D	Time spent in page wait (TOD units)
360	864	PGWAITPG	DS	1D	Time spent in page wait, x pages waiting
368	872	PSASVCCT	DS	1F	Total number of user SVCs
36C	876	PAGELoad	DS	1H	P*1 - Page wait percent, last measurement
36E	878	PAGERATE	DS	1H	P*2 - Paging rate, pages per second
370	880	PSENDCLR	DS	0F	End of area cleared by DMKCPINT
		CPID	DS	1F	CP running identifier
374	884	CPABEND	DS	1F	CP ABEND code
378	888	PSTARTSV	DS	0F	Start of save/restored code
		SYSIPLDV	DS	1H	P*3 - device address of system IPL device
37A	890	PGSRATIO	DC	H'0'	P*4 - Page steals/total replenished
37C	892	ASYSVM	DC	V(DMKSYSVM)	Address of system VMBLOK
380	896	ARSPPR	DC	V(DMKRSPPR)	Address of system printer file chain
384	900	ARSPPU	DC	V(DMKRSPPU)	Address of system punch file chain
388	904	ARSPRD	DC	V(DMKRSPRD)	Address of system reader file chain
38C	908	ARIOPU	DC	V(DMKRIOPU)	Address of system punch table
390	912	ARIOPR	DC	V(DMKRIOPR)	Address of system printer table
394	916	ARIORD	DC	V(DMKRIORD)	Address of system reader table
398	920	PSARV6	DS	1F	Reserved for IBM use
39C	924	ARSPAC	DC	V(DMKRSPAC)	Address of system accounting chain
3A0	928	AVMREAL	DC	A(0)	VMBLOK address of VIRTUAL=REAL user
3A4	932	ASYSABND	DC	A(0)	Address of system ABEND printer
3A8	936	ASYSLC	DC	V(DMKSYSLC)	Address of SYSLOCS information
3AC	940	ASYSOP	DC	V(DMKSYSOP)	Address of system operator VMBLOK
3B0	944	ARIOCT	DC	V(DMKRIOCT)	Address of real channel index table
3B4	948	ARIOCH	DC	V(DMKRIOCH)	Address of first RCHBLOK
3B8	952	ARIOCU	DC	V(DMKRIOCU)	Address of first RCUBLOK
3BC	956	ARIODV	DC	V(DMKRIODV)	Address of first RDEVBLOK
3C0	960	ARIOCC	DC	V(DMKRIOCC)	Address of count of real system channels
3C4	964	ARIOUC	DC	V(DMKRIOUC)	Address of count of real system control units
3C8	968	ARIODC	DC	V(DMKRIODC)	Address of count of real system devices
3CC	972	ACORETBL	DC	V(DMKSYSCS)	Address of system core table
3D0	976	APAGCP	DC	A(X'FFFFFF')	Address of first pageable program
3D4	980	CPCREG0	DC	X'808008C0'	CP architecture control and external mask
3D8	984	CPCREG8	DC	F'0'	Monitor call enable mask
3DC	988	PSARV9	DS	1F	Reserved for IBM use
3E0	992	PSARV10	DS	1F	Reserved for IBM use
3E4	996	PSARV11	DS	1F	Reserved for IBM use
3E8	1000	ADMKFVR	DC	F'0'	Reserved for IBM use
3EC	1004	XVRINST	DC	F'0'	Reserved for IBM use
3F0	1008	PAGECUR	DS	1F	Reserved for IBM use
3F4	1012	MONNEXT	DS	1F	Reserved for IBM use
3F8	1016	PAGEND	DS	1F	Reserved for IBM use
3FC	1020	PAGENXT	DS	1F	Reserved for IBM use
400	1024	TRACEFLG	DS	1F	Trace table flags
404	1028	PSARV12	DS	1F	Reserved for IBM use
408	1032	PSARV15	DS	5D	Reserved for IBM use
430	1072	INSTWRD1	DC	F'0'	Reserved for installation use
434	1076	INSTWRD2	DC	F'0'	Reserved for installation use
438	1080	INSTWRD3	DC	F'0'	Reserved for installation use

43C 1084

INSTWRD4 DC F'0'

Reserved for installation use

Pool of frequently used constants:

440 1088

ZEROES DC 6D'0'

470 1136

BLANKS DC 8X'40'

478 1144

FFS DC 8X'FF'

ALSO = -1

440 1088

F0 EQU ZEROES

480 1152

F1 DC F'1'

484 1156

F2 DC F'2'





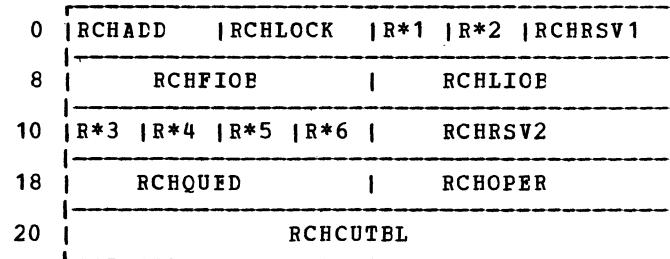
488	1160	F3	DC	F'3'
48C	1164	F4	DC	F'4'
490	1168	F5	DC	F'5'
494	1172	F6	DC	F'6'
498	1176	F7	DC	F'7'
49C	1180	F8	DC	F'8'
4A0	1184	F9	DC	F'9'
4A4	1188	F10	DC	F'10'
4A8	1192	F15	DC	F'15'
4AC	1196	F16	DC	F'16'
4B0	1200	F20	DC	F'20'
4B4	1204	F24	DC	F'24'
4B8	1208	F60	DC	F'60'
4BC	1212	F240	DC	F'240'
4C0	1216	F255	DC	F'255'
4C4	1220	F256	DC	F'256'
4C8	1224	F4095	DC	F'4095'
4CC	1228	F4096	DC	F'4096'
4D0	1232	APTRLK	DC	V(DMKPTRLK)
4D4	1236	NOAED	DC	X'FF000000'
4D8	1240	X4OFFS	DC	X'40FFFFFF'
4DC	1244	XRIGHT24	DC	X'00FFFFFF'
4E0	1248	XPAGNUM	DC	X'00FFF000'
4E4	1252	XRIGHT16	DC	X'0000FFFF'
4E8	1256	AFREE	DC	V(DMKFREE)
4EC	1260	AFRET	DC	V(DMKFRET)
4F0	1264	AQCNWT	DC	V(DMKQCNWT)
4F4	1268	ADSPCH	DC	V(DMKDSPCH)
4F8	1272	APTRAN	DC	V(DMKPTRAN)
4FC	1276	X2048BND	DC	X'00FFF800'
500	1280	PSAEND	DS	0D

ALSO = X'0000000F'

ALSO = X'0000003C'  
 ALSO = X'000000F0' = C'0'  
 ALSO = X'000000FF'  
 ALSO = X'00000100'  
 ALSO = X'00000FFF'  
 ALSO = X'00001000'

End of page 0 usage.

RCHBLOK - REAL CHANNEL BLOCK



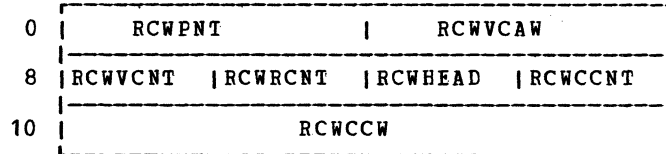
Displacement		Field		Field Description, Contents, Meaning
Hex	Dec	Name		
0	0	RCHADD	DS 1H	Channel address
2	2	RCHLOCK	DS 1H	Channel lock
4	4	RCHSTAT	DS 1X	R*1 - channel status
		Bits Defined in RCHSTAT		
		RCHBUSY	EQU X'80'	Channel busy
		RCHSCED	EQU X'40'	IOB scheduled on channel
		RCHDISA	EQU X'20'	Channel disabled
		RCHDED	EQU X'01'	Channel dedicated
5	5	RCHTYPE	DS 1X	R*2 - Channel type
		Bits Defined in RCHTYPE		
		RCHSEL	EQU X'80'	Selector channel
		RCHBMX	EQU X'40'	Block-multiplexer channel
		RCHIFA	EQU X'81'	Selector-type integrated file adapter
6	6	RCHRSV1	DS 1H	Reserved for IBM use
8	8	RCHFIOB	DS 1F	Pointer to first IOBLOK queued
C	12	RCHLIOB	DS 1F	Pointer to last IOBLOK queued
10	16	RCHDTCK	DS 1X	R*3 - channel data check count
11	17	RCHCCCK	DS 1X	R*4 - channel control check count
12	18	RCHIFCC	DS 1X	R*5 - interface control check count
13	19	RCHCHCK	DS 1X	R*6 - channel chaining check count
14	20	RCHRSV2	DS 1F	Reserved for IBM use
18	24	RCHQUED	DS 1F	IOBLOK queued on channel time
1C	28	RCHOPER	DS 1F	IOBLOK operational on channel time
20	32	RCHCUTBL	DS 32H	Control units attached - RCUSTART index
		RCHSIZE	EQU (*-RCHBLOK)/8	RCHBLOK size in doublewords (X'0C')

RCUBLOK - REAL CONTROL UNIT BLOCK

0	RCUADD	RCULOCK	R*1	R*2	RCURSV1
8	RCUFIOB			RCULIOB	
10	RCUCHA			RCUCHE	
18	RCUQUED			RCUOPER	
20	RCUDVTBL				

Displacement		Field		Field Description, Contents, Meaning
Hex	Dec	Name		
0	0	RCUADD	DS 1H	Control unit address
2	2	RCULOCK	DS 1H	Control unit lock
4	4	RCUSTAT	DS 1X	R*1 - control unit status
		Bits Defined in RCUSTAT		
		RCUBUSY	EQU X'80'	Control unit busy
		RCUSCED	EQU X'40'	IOB scheduled on control unit
		RCUDISA	EQU X'20'	Control unit disabled
		RCUDED	EQU X'01'	Control unit dedicated
5	5	RCUTYPE	DS 1X	R*2 - control unit type
		Bits Defined in RCUTYPE		
		RCUSHRD	EQU X'80'	This control unit can attach to only 1 subchannel
		RCU2701	EQU X'01'	TCU is a 2701
		RCU2702	EQU X'02'	TCU is a 2702
		RCU2703	EQU X'03'	TCU is a 2703
6	6	RCURSV1	DS 1H	Reserved for future use
8	8	RCUFIOB	DS 1F	Pointer to first IOBLOK queued
C	12	RCULIOB	DS 1F	Pointer to last IOBLOK queued
10	16	RCUCHA	DS 1F	Pointer to RCHBLOK - interface A
14	20	RCUCHB	DS 1F	Pointer to RCHBLOK - interface B
18	24	RCUQUED	DS 1F	IOBLOK queued on control unit time
1C	28	RCUOPER	DS 1F	ICBLOK operational on control unit time
20	32	RCUDVTBL	DS 16H	Devices attached - RDVSTART index
		RCUSIZE	EQU (*-RCUBLOK)/8	RCUBLOK size in doublewords (X'08')

RCWTASK - TRANSLATED VIRTUAL I/O CCW



Displacement		Field		Field Description, Contents, Meaning
Hex	Dec	Name		
0	0	RCWPNT	DS 1F	Pointer to next RCWTASK
4	4	RCWVCAW	DS 1F	Virtual address of CCW chain
8	8	RCWVCNT	DS 1H	Virtual CCW count
A	10	RCWRCNT	DS 1H	Real CCW count
C	12	RCWHEAD	DS 1H	RCWTASK header mark X'FFFF'
E	14	RCWCCNT	DS 1H	RCWTASK size in doublewords
10	16	RCWCCW	DS 1D	Cne or more CCWs for device I/O
			ORG RCWCCW	
10	16	RCWADDR	DS 1F	CCW data address
14	20	RCWFLAG	DS 1X	CCW flag bits
15	21	RCWCTL	DS 1X	CCW CP control bits
		Bits Defined in RCWCTL		
		RCWIO	EQU X'80'	I/O data page locked
		RCWGEN	EQU X'40'	CP generated CCW
		RCWHMR	EQU X'20'	DMKUNT to relocate home address/record R0
		RCWREL	EQU X'10'	CCW address relocatable if CCWs moved
		RCWISAM	EQU X'08'	ISAM modifying CCW
		RCW2311	EQU X'04'	TYPE2311T-B pseudo 2311 on 2314
		RCWIDA	EQU X'02'	CP generated indirect data address
16	22	RCWCNT	DS 1H	CCW byte count
			ORG RCWADDR	
10	16	RCWCOMND	DS 1X	CCW command code

RDEVBLK - REAL DEVICE BLOCK

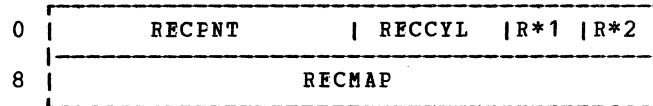
0	RDEVADD	RDEVLOCK	R*1	R*2	R*3	R*4
8	RDEVFIQB		RDEVLIQB			
10	RDEVCUA		RDEVQUE			
18						
20	RDEVIOCT		RDEVAIOB			
28	RDEVUSER	RDEVATT	RDEVCYL			
30		RDEVSER		RDEVLNKS		
38	R	D	E	V	T	C
40	RDEVTMAT		R*5V	R*6	R*7	R*8
48	RDEVICER		RDEVCTRS			

Displacement Hex Dec	Field Name		Field Description, Contents, Meaning
0 0	RDEVADD DS	1H	Device address
2 2	RDEVLOCK DS	1H	Device lock
4 4	RDEVSTAT DS	1X	R*1 - Device status
	Bits Defined in RDEVSTAT		
	RDEVBUSY EQU	X'80'	Device busy
	RDEVSCED EQU	X'40'	IOB scheduled on device
	RDEVDISA EQU	X'20'	Device disabled (offline)
	RDEVRSVD EQU	X'10'	Device reserved
	RDEVIRM EQU	X'08'	Device in intensive error recording mode
	RDEVNRDY EQU	X'04'	Device intervention required
	RDEVDEC EQU	X'01'	Dedicated device (attached to a user)
5 5	RDEVFLAG DS	1X	R*2 - device flags, device dependent
	Bits Defined in RDEVFLAG		
	RDEVSKUP EQU	X'80'	DASD - ascending order seek queuing
	RDEVPREF EQU	X'40'	DASD - volume preferred for paging
	RDEVSYS EQU	X'20'	DASD - volume attached to system
	RDEVOWN EQU	X'10'	DASD - CP owned volume
	RDEVMOUT EQU	X'08'	DASD - volume mounted, not attached
	RDEVPSUP EQU	X'80'	CCNSOLE - terminal has print suppress
	RDEVPREP EQU	X'40'	CCNSOLE - terminal executing prepare command

		RDEVACTV EQU	X'20'	Console - IOBLOK pending; queue request
		RDEVIDNT EQU	X'10'	Console - 2741 terminal code identified
		RDEVENAB EQU	X'08'	Console - device is enabled
		RDEVHIO EQU	X'04'	Console - next interrupt from a halt I/O
		RDEVDISB EQU	X'02'	Console - device is to be disabled
		RDEVDRAN EQU	X'80'	Spooling - device output drained
		RDEVTERM EQU	X'40'	Spooling - device output terminated
		RDEVACNT EQU	X'20'	Spooling - device busy with accounting
		RDEVSPAC EQU	X'10'	Spooling - force printer to single space
		RDEVRSR EQU	X'08'	Spooling - restart current file
		RDEVBACK EQU	X'04'	Spooling - backspace the current file
		RDEVSEP EQU	X'02'	Spooling - print/punch job separator
		RDEVLOAD EQU	X'01'	Spooling - UCS buffer verified
6	6	RDEVTPC DS	1X	R*3 - device type class (See Appendix C)
7	7	RDEVTYPE DS	1X	R*4 - device type (See Appendix C)
8	8	RDEVFICB DS	1F	Pointer to first IOBLOK queued
C	12	RDEVLICB DS	1F	Pointer to last IOBLOK queued
10	16	RDEVCUA DS	1F	Pointer to RCUBLOK - interface A
14	20	RDEVSUB DS	1F	Pointer to RCUBLOK - interface B
18	24	RDEVQUED DS	1D	IOBLOK queued time - TOD clock units
20	32	RDEVIOCT DS	1F	Device I/O count
24	36	RDEVAICB DS	1F	Active IOBLOK
28	40	RDEVUSER DS	1F	Pointer to VMBLOK of dedicated user
2C	44	RDEVATT DS	1H	Attached virtual address
2E	46	RDEVCYL DS	1H	DASD - current cylinder location
30	48	RDEVSER DS	CL6	Device volume serial number
36	54	RDEVLNKS DS	1H	DASD - number of links to this disk
38	56	RDEVTCTL DS	8X	Terminal control bytes
40	64	RDEVTMAT DS	1F	Device attached time - TOD clock word 0
44	68	RDEVRSV1 DS	1X	R*5 - Reserved for IBM use
45	69	RDEVSTA2 DS	1X	R*6 - Device status (2nd byte)
		Bits Defined in RDEVSTA2		
		RDEVRACT EQU	X'80'	Active device is being reset
		RDEVBUCH EQU	X'40'	Device is busy with the channel
46	70	RDEVMDL DS	1X	R*7 - device model number
47	71	RDEVFTR DS	1X	R*8 - device feature code
48	72	RDEVIOFR DS	1F	Pointer to IOERBLOK for last CP error
4C	76	RDEVCTRS DS	1F	Pointer to error counter control blok
		RDEVSIZE EQU	(*-RDEVBLOK)/8	RDEVBLOK size in doublewords (X'0A')
		For CP owned devices		
		ORG	RDEVUSER	
28	40	RDEVALLN DS	1F	Anchor for ALOCBLOK chain for this device
2C	44	RDEVCODE DS	1H	Device code - SYSOWNED index
		ORG	RDEVTCTL	
38	56	RDEVPAGE DS	1F	Anchor for RECBLOK chain for paging
3C	60	RDEVRECS DS	1F	Anchor for RECBLOK chain for spooling

40	64	RDEVPNT DS	1F	Pointer to next RDEVBLK for allocation
		For slotted 2301 paging devices		
		ORG RDEVRECS		
3C	60	RDEVDCIL DS	1F	Pointer to DRUMTABL control block
		For spooling unit record devices		
		ORG RDEVQUED		
18	24	RDEVSPIL DS	1F	Pointer to active RSPLCTL block
1C	28	RDEVCLAS DS	4C	Device class(es)
		For terminal devices		
		ORG RDEVQUED		
18	24	RDEVCON DS	1F	Pointer to CONTASK list
1C	28	RDEVAIRA DS	1F	Attention interrupt return address
		ORG RDEVTCTL		
38	56	RDEVLEND DS	1C	Device line end symbol
39	57	RDEVLDEL DS	1C	Device line delete symbol
3A	58	RDEVCDL DS	1C	Device character delete symbol
3B	59	RDEVESCP DS	1C	Device character escape symbol
3C	60	RDEVLEN DS	1X	Device line length
3D	61	RDEVATUC DS	1X	Device attention count
3E	62	RDEVTFLG DS	1X	Additional terminal flags
		Bits Defined in RDEVTFLG		
		RDEVATIN EQU	X'80'	Attention signalled on input
		RDEVREST EQU	X'40'	Terminal in process of being reset
		RDEVATCF EQU	X'20'	Do not type exclamation point or CR
		RDEVCIIRD EQU	X'10'	Write a circle D to a terminal
3F	63	RDEVRSV3 DS	1X	Reserved for IBM use
		ORG RDEVMDL		
46	70	RDEVTMCD DS	1X	Terminal code
		Bits Defined in RDEVTMCD		
		RDEVPTIC EQU	X'00'	PTTC/EBCD
		RDEVCOFR EQU	X'04'	Correspondence
		RDEVAPLP EQU	X'08'	APL PTTC/EBCD
		RDEVAPLC EQU	X'0C'	APL Correspondence
		RDEVUSC8 EQU	X'10'	UASCII-8 level
47	71	RDEVSDN DS	1X	Terminal set-address number

RECBLOCK - DASD PAGE (SLOT) ALLOCATION BLOCK



Displacement Hex Dec	Field Name	Field Description, Contents, Meaning
0 0	RECPNT DS 1F	Pointer to next RECBLOCK on chain
4 4	RECCYL DS 1H	Cylinder address for pages in this block
6 6	RECUSED DS 1X	R*1 - Number of pages currently in use
7 7	RECMAX DS 1X	R*2 - Maximum number of pages available
8 8	RECMAP DS 1D	Page allocation bit map
	Bits Defined in RECMAP	
	0 - Page is available	
	1 - Page has been assigned	
	RECSIZE EQU (*-RECBLOCK)/8 RECBLOCK size in doublewords (X'02')	

Note: Although the size of RECMAP is fixed, the maximum number of pages available on a cylinder is device dependent. For any pages that are not physically present on a cylinder, their corresponding bits are set to one.



RSPCTL - REAL SPOOL CONTROL BLOCK

0	RSPRSIRT		RSPDPAGE
8	RSPVPAGE		RSPRPAGE
10	RSPMISC		RSPSFELK

Displacement		Field			Field Description, Contents, Meaning
Hex	Dec	Name			
0	0	RSPRSTRT DS	1F		Restart CAW - CCW address
4	4	RSPDPAGE DS	1F		DASD location (DCHR) of current page buffer
8	8	RSPVPAGE DS	1F		Virtual address of page buffer
C	12	RSPRPAGE DS	1F		Real address of page buffer
10	16	RSPMISC DS	1F		Use varies according to caller
14	20	RSPSFBLK DS	1F		Pointer to SFBLOK for file
		RSPSIZE EQU	(*-RSPCTL)/8		Size in doublewords (X'03')

SAVEAREA

0	SAVERETN		SAVER12
8	SAVER13		SAVEWRK1
10	SAVEREGS		
40	SAVEWRK2		SAVEWRK3
48	SAVEWRK4		SAVEWRK5
50	SAVEWRK6		SAVEWRK7
58	SAVEWRK8		SAVEWRK9

Displacement		Field		Field Description, Contents, Meaning
Hex	Dec	Name		
0	0	SAVERETN	DS 1F	Active SAVEAREA (caller's return address)
			ORG SAVERETN	
0	0	SAVENEXT	DS 1F	Inactive SAVEAREA (next SAVEAREA address)
4	4	SAVER12	DS 1F	Caller's base (R12)
8	8	SAVER13	DS 1F	Caller's SAVEAREA (R13)
C	12	SAVEWRK1	DS 1F	Callee's workarea
10	16	SAVEREGS	DS 12F	Caller's registers (R0 TO R11)
			ORG SAVEREGS	
10	16	SAVER0	DS 1F	
14	20	SAVER1	DS 1F	
18	24	SAVER2	DS 1F	
1C	28	SAVER3	DS 1F	
20	32	SAVER4	DS 1F	
24	36	SAVER5	DS 1F	
28	40	SAVER6	DS 1F	
2C	44	SAVER7	DS 1F	
30	48	SAVER8	DS 1F	
34	52	SAVER9	DS 1F	
38	56	SAVER10	DS 1F	
3C	60	SAVER11	DS 1F	
40	64	SAVEWRK2	DS 1F	Callee's workarea (8 words)
44	68	SAVEWRK3	DS 1F	
48	72	SAVEWRK4	DS 1F	
4C	76	SAVEWRK5	DS 1F	

54 84  
58 88  
5C 92

SAVEWRK7 DS 1F  
SAVEWRK8 DS 1F  
SAVEWRK9 DS 1F

SAVESIZE EQU (\*-SAVEAREA)/8 size in doublewords (X'0C')

SAVTABLE - FIRST PAGE ON SAVED SYSTEM DASD

0	SAVPSW
8	SAVGREGS
48	SAVFPRES
68	SAVCREGS
A8	SAVKEYS

Displacement		Field	Field Description, Contents, Meaning		
Hex	Dec	Name			
0	0	SAVPSW DS 1D			PSW of virtual machine at SAVSYS table
8	8	SAVGREGS DS 16F			General registers
48	72	SAVFPRES DS 4D			Floating-point registers
68	104	SAVCREGS DS 16F			Control registers
A8	168	SAVKEYS DS 1H			Two byte entry for each saved page containing storage keys for each page

SEGTABLE - SEGMENT TABLE

0	S*1 SEGPAGE (variable length)
---	----------------------------------

Displacement		Field	Field Description, Contents, Meaning		
Hex	Dec	Name			
0	0	SEGPAGE DS 1F			Pointer to page table - PAGTABLE
		Page table length			
		ORG SEGPAGE			
0	0	SEGPLN DS 1X			S*1 - page table length (pages - 1) (in left half of byte)

SFBLOK - SPOOL FILE BLOCK

0	SFBPNT		SFBSTART
8	SFBUSER		
10	SFBORIG		
18	SFBRECNO	SFBRECSZ	SFBFILID
20	S*1	S*2	SFBMISC1   SFBRECS
28	SFBFNAME		
34	SFBFTYPE		
40	SFBDATE		
48	SFBTIME		
50	SFELAST	SFBCOPY	S*3  S*4
58	SFBDIST		

Displacement  
Hex Dec

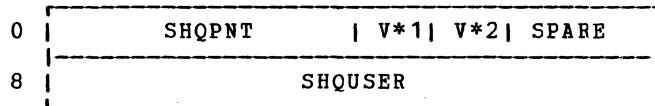
Field  
Name

Field Description, Contents, Meaning

0	0	SFBPNT	DS	1F	Pcinter to next SFBLOK
4	4	SFBSTART	DS	1F	DASD location (DCHR) of last page buffer
8	8	SFBUSER	DS	CL8	VMUSER identification of file owner
10	16	SFBORIG	DS	CL8	VMUSER identification of file origin
18	24	SFBRECNO	DS	1F	Number of data records in file
1C	28	SFBRECSZ	DS	1H	Logical record size - excluding CCW's
1E	30	SFBFILID	DS	1H	Binary system file number
20	32	SFBFLAG	DS	1X	S*1 - SFBLOK control bits
Bits Defined in SFBFLAG					
		SFBINUSE	EQU	X'80'	File being processed
		SFBRECCK	EQU	X'40'	Allocation records complete
		SFBUHOLD	EQU	X'20'	File in user hold status
		SFBDUMP	EQU	X'10'	File is a CP system dump
		SFBOPEN	EQU	X'08'	Input file has been opened
		SFBSHOLD	EQU	X'04'	File in system hold status
		SFBEOF	EQU	X'02'	Input file has reached EOF
		SFBRECER	EQU	X'01'	SFBREC chain incomplete
21	33	SFBTYPE	DS	1X	S*2 - device type for output

22	34	SFBMISC1 DS	1H	Use varies according to caller	0
24	36	SFBRECS DS	1F	Pointer to RECBLOKS for active file	0
28	40	SFBFNAME DS	CL12	File name	
34	52	SFBFTYPE DS	CL12	File type	
40	64	SFBDATE DS	CL8	Creation date of spool file	MM/DD/YY
48	72	SFBTIME DS	CL8	Creation time of spool file	MM:SS
50	80	SFBLAST DS	1F	DASD location (DCHR) of last page buffer	D1100000
54	84	SFBCOPY DS	1H	Number of copies requested	
56	86	SFBCLAS DS	1C	S*3 - Spool output class	
57	87	SFBFLAG2 DS	1X	S*4 - SFBLOK flag byte two	
Bits Defined in SFBFLAG2					
		SFBHOLD EQU	X'80'	Save input file, or hold output file	
		SFBNOHID EQU	X'40'	Delete input file, or do not hold output file	
		SFBHOLD and SFBNOHID		Override options in VDEVBLK	
		SFBREQUE EQU	X'20'	Re-queue spool file	
		SFBRSTRT EQU	X'10'	Restart in progress	
		SFBTICER EQU	X'08'	Buffer TIC error	
		SFBPURGE EQU	X'04'	Purge open spool file	
58	88	SFBDIST DS	CL8	Distribution code	
		SFBSIZE EQU	(*SFBLOK)/8	Size in doublewords (X'0C')	00000000 0000100

SHQBLOK - SPOOL HOLD QUEUE BLOCK



Displacement	Field	Field	Field	Field
Hex Dec	Name			Description, Contents, Meaning
0 0	SHQPNT DS	1F		Address of next SHQBLOK
4 4	SHQFLAGS DS	0CL4		Length
4 4	SHQUHOLD DS	1X		V*1 user 'USER HOLD' flag byte
5 5	SHQSHOLD DS	1X		V*2 user 'SYSTEM HOLD' flag byte
6 6	SHQSPARE DS	2X		Spare
8 8	SHQUSER DS	CL8		VMUSER identification of file owner

SHQBSIZE EQU (\*SHQBLOK)/8 Size in doublewords (X'02')

Bits Defined in SHQUHOLD and SHQSHOLD  
 TYPRT is used for printer type. (See Appendix C for DEVTPES)  
 TYFPUN is used for punch type. (See Appendix C for DEVTPES)

SHRTABLE - NAMED-SHARED SEGMENT SYSTEMS

0	SHRFPNT		SHRBPNT
8	SHRNAME		
10	SHRTSIZE		SHRSEGCT
18	SHRSEGNM		SHRPAGE

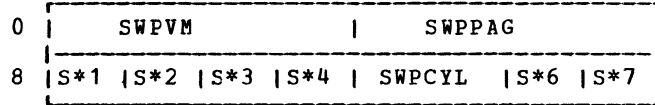
Displacement Hex Dec	Field Name			Field Description, Contents, Meaning
0 0	SHRFPNT DS	1F		Pointer to next SHRTABLE
4 4	SHRBPNT DS	1F		Pointer to previous SHRTABLE
8 8	SHRNAME DS	CL8		Name of saved system
10 16	SHRTSIZE DS	1H		Size of SHRTABLE in doublewords
12 18	SHRUSECT DS	1H		Number of users IPLed to this name
14 20	SHRSEGCT DS	1F		Number of shared segments
18 24	SHRSEGNM DS	1F		Contains shared segment numbers. Up to four segment numbers per word.
1C 28	SHRPAGE DS	1F		Pointers to each of the shared SEGTABLES. There is one word for each shared segment. The entry is the same as S*1 SEGPAGE in the SEGTABLE.

SPLINK - SPOOL PAGE BUFFER LINKAGE

0	SPNXTFAG		SPPREPAG
8	SPRMISC		SPRECNUM
10	Spool Buffer Data		

Displacement Hex Dec	Field Name			Field Description, Contents, Meaning
0 0	SPNXTFAG DS	1F		DASD location (DCHR) of next page buffer
4 4	SPPREPAG DS	1F		DASD location (DCHR) of previous page buffer
8 8	SPRMISC DS	1F		Use varies according to caller
C 12	SPRECNUM DS	1F		Number of data records in buffer
	SPSIZE EQU	(*-SPLINK)		Size in bytes (X'10')

SWPTABLE - SWAP TABLE FOR VIRTUAL MACHINE PAGING



Displacement		Field			Field Description, Contents, Meaning
Hex	Dec	Name			
0	0	SWPVM	DS	1F	Pointer to VMBLOK
4	4	SWPPAG	DS	1F	Pointer to PAGTABLE
8	8	SWPFLAG	DS	1X	S*1 - SWPTABLE flag bits
		Bits Defined in SWPFLAG			
		SWPTRANS	EQU	X'80'	Page in transit
		SWPRECMP	EQU	X'40'	Page permanently assigned
		SWPALLCC	EQU	X'20'	Page enqueued for allocation
		SWPSHR	EQU	X'10'	Page shared
		SWPREF1	EQU	X'08'	First half page referenced
		SWPCHG1	EQU	X'04'	First half page changed
		SWPREF2	EQU	X'02'	Second half page referenced
		SWPCHG2	EQU	X'01'	Second half page changed
9	9	SWVPV PAGE	DS	1X	S*2 - virtual page number
A	10	SWPKEY1	DS	1X	S*3 - virtual storage key
B	11	SWPKEY2	DS	1X	S*4 - virtual storage key
C	12	SWPCYL	DS	1H	DASD cylinder address
E	14	SWPDPAGE	DS	1X	S*6 - Page number on cylinder
F	15	SWPCODE	DS	1X	S*7 - RDEVBLK device code

Note: For each SWPTABLE there is only one doubleword that consists of SWPVM and SWPPAG followed by 16 entries (one for each PAGTABLE entry) that consist of S\*1, S\*2, S\*3, S\*4, SWPCYL, S\*6 and S\*7. Thus, the total size of the SWPTABLE is 17 doublewords.

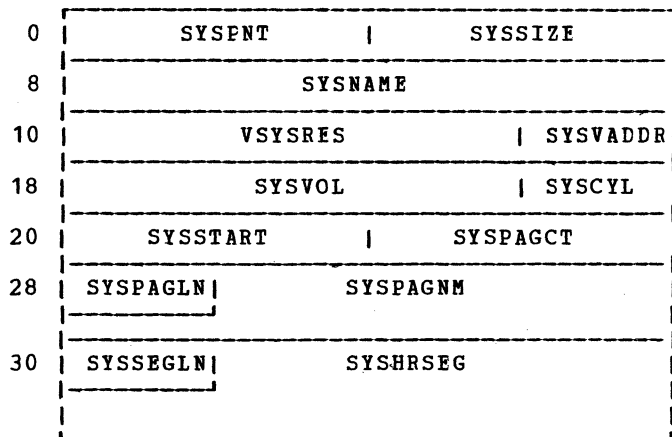


SYSLOCS - SYSTEM LCW STORAGE INFORMATION

0	DMKSYSDT	DMKSYSTEM
10	DMKSYSLW	DMKSYSLG
20	DMKSYSNM	DMKSYSMA DMKSYSMU  DMKSYSND
30	DMKSYSLB	DMKSYSUD DMKSYSPL
40	DMKSYSDW	#  #  @  "
50	S*1	DMKSYSCK

Displacement		Field		Field Description, Contents, Meaning
Hex	Dec	Name		
0	0	DMKSYSDT DC	CL8'MM/DD/YY'	Date of system log message
8	8	DMKSYSTEM DC	CL8'HH:MM:SS'	Time of system log message
10	16	DMKSYSLW DC	X'00',X'00',CL10'	' Weekday of system log message
1C	28	DMKSYSLG DC	A(0)	Pointer to first log message block
20	32	DMKSYSNM DC	F'0'	Current number of users on the system
24	36	DMKSYSMA DC	F'0'	Maximum number of users allowed on
28	40	DMKSYSMU DC	F'0'	Maximum number of users on the system
2C	44	DMKSYSND DC	F'0'	Number of dialed users on the system
30	48	DMKSYSLB DC	A(0)	Pointer to user directory lock block
34	52	DMKSYSUD DC	A(0)	Pointer to start of user directory on SYSRES
38	56	DMKSYSPL DC	A(0)	Pointer to a list of virtual page buffers
3C	60	DC	A(0)	Reserved for IBM use
40	64	DMKSYSDW DC	X'00',X'00',CL10'	' Day-of-week in Hex. and EBCDIC
4C	76	DMKSYSLE DC	X'7B'	# default line-end (pound-sign)
4D	77	DMKSYSLE DC	X'4A'	ø default line-delete (cent-sign)
4E	78	DMKSYSLE DC	X'7C'	@ default character-delete (at-sign)
4F	79	DMKSYSLE DC	X'7F'	" default edit escape (double-quote-mark)
50	80	DMKSYSLL DC	AL1(130,129,72)	S*1 Default line lengths for 3210 & 3215 - 2741 & 1050 - TTY terminals
53	83	DC	XL5'0'	Reserved for IBM use
58	88	DMKSYSCK DC	D'0'	Time-of-day clock value last stored by accounting, DUMP or machine check

SYSTABLE - NAMED SYSTEM TABLE



Displacement		Field			Field Description, Contents, Meaning
Hex	Dec	Name			
0	0	SYSPNT	DS	1F	Chain pointer to next entry
4	4	SYSSIZE	DS	1F	Minimum storage size needed to run system
8	8	SYSNAME	DS	CL8	System name
10	16	VSYSRES	DS	CL6	Volume serial of DASD containing user's system
16	22	YSVADDR	DS	1H	Virtual address of VSYSRES
18	24	SYSVOL	DS	CL6	Volume serial of DASD containing saved pages
1E	30	SYSCYL	DS	1H	Cylinder on VSYSRES of user's system same as VDEVRELN
20	32	SYSSTART	DS	1F	CCPD of first page on SYSVOL
24	36	SYSPAGCT	DS	1F	Total number of pages saved
28	40	SYSPAGLN	DS	1H	Number of entries in SYSPAGNM
2C	44	SYSPAGNM	DS	1F	One full word entry for each range of pages to be saved
30	48	SYSSEGLN	DS	1H	Numbers of entries in SYSHRSEG
32	50	SYSHRSEG	DS	1X	One byte for each segment to be shared

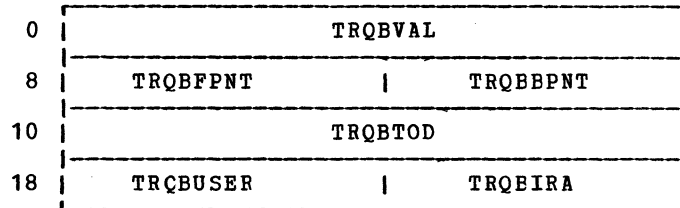
TREXT - VIRTUAL MACHINE TRACING EXTENSION TO VMELOK

0	TREXIN1		TREXIN2
8	TREXSVC1	TREXSVC2	S*1   S*2   TREXLOCK
10	TREXPERA		TREXPERC   TREXLCNT
18	TREXANSI		TREXCR9
20	TREXCR10		TREXCR11
28	TREXBUFF		

Displacement Hex Dec	Field Name			Field Description, Contents, Meaning
0 0	TREXIN1 DS	1F		First address - replaced instruction
4 4	TREXIN2 DS	1F		Second address - replaced instruction
8 8	TREXSVC1 DS	1H		Displaced halfword - instruction one
A 10	TREXSVC2 DS	1H		Displaced halfword - instruction two
		ORG	TREXIN1	
0 0	TREXPSW DS	1D		Old PSW for pending SVC interrupt
8 8	TREXINTL DS	1H		Instruction length code
A 10	TREXINTC DS	1H		Interruption code for pending interrupt
C 12	TREXFLAG DS	1X		S*1 - tracing control flags
	Bits Defined in TREXFLAG			
	TREXRUN EQU	X'80'		Prevent CFWAIT between events
	TREXVAT EQU	X'40'		Call DMKVATR to put back virtual instruction
D 13	TREXOUT DS	1X		S*2 - trace output controls
	Bits Defined in TREXOUT			
	TREXPRT EQU	X'80'		Output to the virtual printer
	TREXCON EQU	X'40'		Output to user terminal
E 14	TREXLOCK DS	1H		Indicates tracing when set
10 16	TREXPERA DS	1F		PER event address on interrupt
14 20	TREXPERC DS	1H		PER code bits from hardware event
16 22	TREXLCNT DS	1H		Printed output line count
18 24	TREXANSI DS	1A		Address of next (or last) sequential instruction
1C 28	TREXCR9 DS	0F		Shadow control registers for PER trace
	TREXPER DS	XL2		PER control field
1E 30	TREXPREG DS	1H		PER register mask field
20 32	TREXCR10 DS	1F		Address range start value
24 36	TREXCR11 DS	1F		Address range ending value

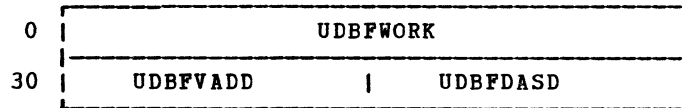
28	40	TREXBUFF DS	10D	Console/printer output buffer (80 bytes)
		TREXSIZE EQU	(*-TREXT)/8	TREXT size in doublewords (X'0F')
10	16	TREXNSI DS	6X	Actual next (or last) sequential instruction
		ORG	TREXPERA	Re-Definition for TRACE use
1C	28	TREXCTL DS	0H	Halfword holding tracing control bits:
		TREXCTL1 DS	1X	First byte = same as VMTRCTL in VMBLOK
1D	29	TREXCTL2 DS	1X	Second byte = remaining control bits
		Bits Defined in TREXCTL2:		
		TREXCCW EQU	X'80'	TRACE virtual and real CCWs
		TREXCSW EQU	X'40'	TRACE virtual and real CSWs
		TREXBRAN EQU	X'20'	TRACE successful branches
		TREXINST EQU	X'10'	TRACE all instructions
1E	30	TREXPRNT DS	1H	Printer flagbits corresponding to TREXCTL
20	32	TREXTERM DS	1H	Terminal flagbits corresponding to TREXCTL
22	43	TREXRUNF DS	1H	Run/norun flagbits corresponding to TREXCTL
24	45	TREXPNTR DS	1F	Pointer to 1st stacked TRACE request, if any

TROBLOCK - TIMER REQUEST FLOCK



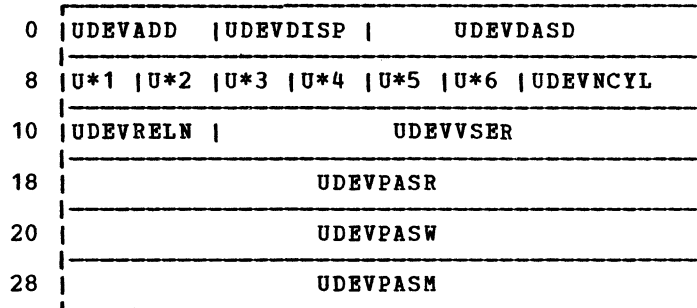
Displacement	Field	Field Description, Contents, Meaning
Hex Dec	Name	
0 0	TROBVAL DS 1D	TCD clock comparator value for interrupt
8 8	TROBFPNT DS 1F	Pointer to next TROBLOK
C 12	TROBBPNT DS 1F	Pointer to previous TROBLOK
10 16	TROBTOD DS 1D	TCD clock value when TROBLOK is queued
18 24	TROBUSER DS 1F	Address of VMBLOK for user
1C 28	TROBIRA DS 1F	Interrupt return address
	TROBSIZE EQU (*-TROBLOK)/8	Size in doublewords (X'04')

UDBFBLOK - USER DIRECTORY BUFFER BLOCK



Displacement Hex Dec	Field Name	Field Description, Contents, Meaning
0 0	UDBFWORK DS 6D	Buffer work space used by the caller
30 48	UDBFVADD DS 1F	Virtual address of the last directory page
34 52	UDBFDASD DS 1F	DASD address of the last directory page
	UDBFSIZE EQU (*-UDBFBLOK)/8	UDBFBLOK size in doublewords (X'07')

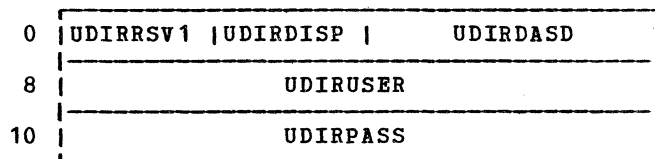
UDEVBLOK - USER DEVICE BLOCK



Displacement Hex Dec	Field Name	Field Description, Contents, Meaning
0 0	UDEVADD DS 1H	Virtual device address
2 2	UDEVDISP DS 1H	Displacement of the next UDEVBLOK
4 4	UDEVNASD DS 1F	DASD address of the next UDEVBLOK
8 8	UDEVSTAT DS 1X	U*1 - status information
	Bits Defined in UDEVSTAT	
	UDEVDED EQU X'80'	Device to be dedicated to this user
	UDEVTDISK EQU X'40'	TDISK to be allocated
	UDEVLONG EQU X'20'	Device block is full length (6 doublewords)
	UDEVLKDV EQU X'10'	Device is to be linked (at LOGON time)
	UDEVSPOO EQU X'08'	Device is a spool device

9	9	UDEVMODE DS	1X	U*2 - access mode information
		Bits Defined in UDEVMODE		
		UDEVLR EQU	X'80'	Read links allowed
		UDEVLW EQU	X'40'	Write links allowed
		UDEVLM EQU	X'20'	Multiple-write links allowed
		UDEVVR EQU	0	Device to be in R link mode for owner
		UDEVRR EQU	4	Device to be in RR link mode for owner
		UDEVW EQU	8	Device to be in W link mode for owner
		UDEVWR EQU	12	Device to be in WR link mode for owner
		UDEVV EQU	16	Device to be in M link mode for owner
		UDEVMR EQU	20	Device to be in MR link mode for owner
		UDEVMW EQU	24	Device to be in MW link mode for owner
A	10	UDEVTPC DS	1C	U*3 - device class
B	11	UDEVTYPE DS	1C	U*4 - device type
C	12	UDEVFTR DS	1C	U*5 - device feature mode
D	13	UDEVMDL DS	1C	U*6 - device model number
E	14	UDEVNCYL DS	1H	Virtual DASD size
10	16	UDEVRELN DS	1H	Virtual DASD cylinder relocation
12	18	UDEVVSER DS	6C	Volume serial number
18	24	UDEVPASR DS	1D	Password for read access
20	32	UDEVPASW DS	1D	Password for write access
28	40	UDEVPASM DS	1D	Password for multiple access
		UDEVSIZE EQU	(*-UDEVBLOK)/8	UDEVBLOK size in doublewords
		ORG	UDEVMDL	
D	13	UDEVCLAS DS	1C	User device block (Short)
E	14	UDEVLINK DS	1H	U*6 - unit spool output class
10	16	UDEVLKID DS	1D	User link to USERID

UDIRBLOK - USER DIRECTORY BLOCK

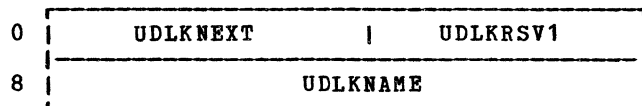


Displacement			Field		Field Description, Contents, Meaning
Hex	Dec		Name		
0	0		UDIRRSV1 DS	1H	Reserved for IBM use
2	2		UDIRDISP DS	1H	Displacement of the user UMACBLOK

4	4	UDIRDASD DS	1F	DASD address of the user UMACBLOK
8	8	UDIRUSER DS	1D	USERID
10	16	UDIRPASS DS	1D	User password

UDIRSIZE EQU (\*-UDIRBLOK)/8 UDIRBLOK size in doublewords (X'03')

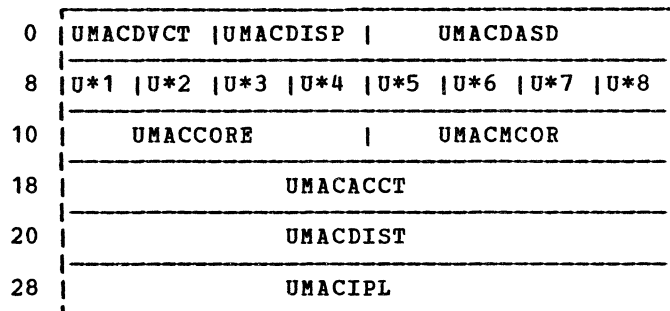
UDLKBLOK - USER DIRECTORY LOCK BLOCK



Displacement Hex Dec	Field Name		Field Description, Contents, Meaning
0 0	UDLKNEXT DS	1F	Pointer to the next lock block
4 4	UDLKRSV1 DS	1F	Reserved for IBM use
8 8	UDLKNAME DS	1D	The name locked

UDLKSIZE EQU (\*-UDLKBLOK)/8 UDLKBLOK size in doublewords (X'02')

UMACBLOK - USER MACHINE BLOCK

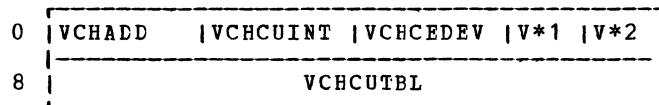


Displacement Hex Dec	Field Name		Field Description, Contents, Meaning
0 0	UMACDVCT DS	1H	Number of devices
2 2	UMACDISP DS	1H	Displacement of the user's first UDEVBLOK

4	4	UMACDASD DS	1F	DASD address of the user's first UDEVBLK
8	8	UMACCLEV DS	1C	U*1 - command level
		Bits Defined in UMACCLEV		
		UMACCLA EQU	X'80'	Class A functions
		UMACCLB EQU	X'40'	Class B functions
		UMACCLC EQU	X'20'	Class C functions
		UMACCLD EQU	X'10'	Class D functions
		UMACCLE EQU	X'08'	Class E functions
		UMACCLF EQU	X'04'	Class F functions
		UMACCLG EQU	X'02'	Class G functions
		UMACCLH EQU	X'01'	Class H functions
9	9	UMACPRIR DS	1X	U*2 - priority
A	10	UMACOPT DS	1X	U*3 - virtual machine options
		Bits Defined in UMACOPT		
		UMACISAM EQU	X'80'	ISAM CCW checking option
		UMACECOP EQU	X'40'	Extended control mode option
		UMACRT EQU	X'20'	Real timer option
		UMACVROP EQU	X'10'	Virtual = Real storage option
		UMACACC EQU	X'08'	Accounting card option
B	11	UMACRSV1 DS	1C	U*4 - Reserved for IBM use
C	12	UMACLEND DS	1C	U*5 - Terminal line end symbol
D	13	UMACLDEL DS	1C	U*6 - Terminal line delete symbol
E	14	UMACCDEL DS	1C	U*7 - Terminal character delete symbol
F	15	UMACES DS	1C	U*8 - Edit escape symbol
10	16	UMACCORE DS	1F	Virtual storage size in bytes
14	20	UMACMCOR DS	1F	Maximum virtual storage size in bytes
18	24	UMACACCT DS	1D	Accounting information
20	32	UMACDIST DS	1D	User machine distribution information
28	40	UMACIPL DS	1D	Name of system to be IPLed at LOGON time
		UMACSIZE EQU	(*-UMACBLOK)/8	UMACBLOK size in doublewords (X'06')

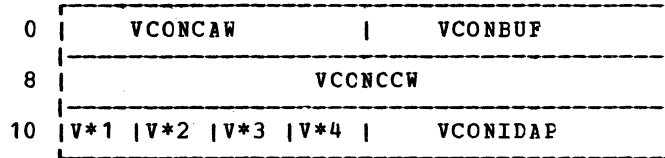


VCHBLOK - VIRTUAL CHANNEL BLOCK



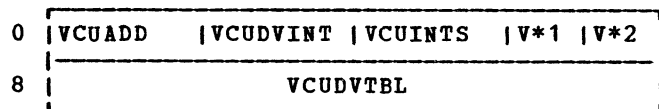
Displacement		Field		Field Description, Contents, Meaning
Hex	Dec	Name		
0	0	VCHADD	DS 1H	Virtual channel address
2	2	VCHCUINT	DS 1H	VCUBLOCK with interrupt - bit map
4	4	VCHCEDEV	DS 1H	Virtual device address with channel class interrupt
6	6	VCHSTAT	DS 1X	V*1 - virtual channel status
		Bits Defined in VCHSTAT		
		VCHBUSY	EQU X'80'	Virtual channel busy
		VCHCEPND	EQU X'40'	Virtual channel class interrupt pending
		VCHDED	EQU X'01'	Virtual channel dedicated
7	7	VCHTYPE	DS 1X	V*2 - Virtual channel type
		Bits Defined in VCHTYPE		
		VCHSEL	EQU X'80'	Virtual selector channel
		VCHBMX	EQU X'40'	Virtual block multiplexor
8	8	VCHCUTBL	DS 16H	Control units attached - VMCUSTRT index
		VCHSIZE	EQU (*-VCHBLOK)/8	VCHBLOK size in doublewords (X'05')

VCONCTL - VIRTUAL CONSOLE CONTROL BLOCK



Displacement		Field		Field Description, Contents, Meaning
Hex	Dec	Name		
0	0	VCONCAW	DS 1F	Virtual address of user CCW
4	4	VCONBUF	DS 1F	Pointer to data buffer
8	8	VCONCCW	DS 1D	Current user CCW
10	16	VCONRSV1	DS 1X	V*1 - Reserved for IBM use
11	17	VCONBFSZ	DS 1X	V*2 - Data buffer size in doublewords
12	18	VCONRSV2	DS 1X	V*3 - reserved for IBM use
13	19	VCONRSV3	DS 1X	V*4 - reserved for IBM use
14	20	VCONIDAP	DS 1F	For IDA pointer to current IDAW
		VCONSIZE	EQU (*-VCONCTL)/8	VCONCTL size in doublewords (X'03')
			ORG VCONCCW	
8	8	VCONADDR	DS 1F	CCW data address
C	12	VCONFLAG	DS 1X	CCW flag bits
D	13	VCONRSV4	DS 1X	Reserved for IBM use
E	14	VCONCNT	DS 1H	CCW byte count
			ORG VCONADDR	
8	8	VCONCMD	DS 1X	CCW command code

VCUBLOK - VIRTUAL CONTROL UNIT BLOCK



Displacement		Field		Field Description, Contents, Meaning
Hex	Dec	Name		
0	0	VCUADD	DS 1H	Virtual control unit address
2	2	VCUDVINT	DS 1H	VDEVBLK with interrupt - bit map
4	4	VCUINTS	DS 1H	Virtual control unit interrupt status
6	6	VCUSTAT	DS 1X	V*1 - virtual control unit status
		Bits Defined in VCUSTAT		
		VCUCHBSY	EQU X'80'	Virtual subchannel busy
		VCUCEPND	EQU X'40'	Interrupt pending in sub-channel
		VCUBUSY	EQU X'20'	Virtual control unit busy
		VCUPEND	EQU X'10'	Virtual control unit interrupt pending
		VCUCUEPN	EQU X'08'	Virtual control unit end pending
7	7	VCUTYPE	DS 1X	V*2 - virtual control unit type
		Bits Defined in VCUTYPE		
		VCUSHRD	EQU X'80'	Virtual control unit on shared subchannel
		VCUCTCA	EQU X'40'	Virtual control unit is a channel-to-channel adapter
8	8	VCUDVTBL	DS 16H	Devices attached - VMDVSTRT index
		VCUSIZE	EQU (*-VCUBLOK)/8	VCUBLOK size in doublewords (X'05')

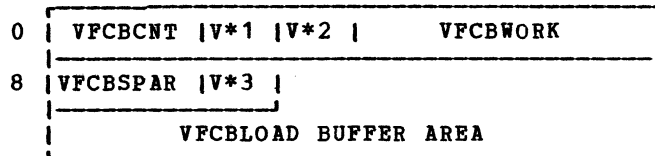
VDEVBLOK - VIRTUAL DEVICE BLOCK

0	VDEVADD	VDEVINTS	V*1	V*2	V*3	V*4
8	VDEVCSW					
10	VDEVRELN	VDEVBND	VDEVPOSN			
18	VDEVQUED	VDEVOPER				
20	VDEVLINK	VDEVEREAL				
28	VDEVIOCT	VDEVUSER				
30	VDEVIOER	VDEVIOB				

Displacement Hex Dec	Field Name	Field Description, Contents, Meaning
0 0	VDEVADD DS 1H	Virtual Device Address
2 2	VDEVINTS DS 1H	Virtual Device Interrupt Status
4 4	VDEVTPC DS 1X	V*1 - virtual device type class
5 5	VDEVTYPE DS 1X	V*2 - virtual device type
6 6	VDEVSTAT DS 1X	V*3 - virtual device status
	Bits Defined in VDEVSTAT	
	VDEVCHBS EQU X'80'	Virtual subchannel busy
	VDEVCHAN EQU X'40'	Virtual channel interrupt pending
	VDEVBUSY EQU X'20'	Virtual device busy
	VDEVPEND EQU X'10'	Virtual device interrupt pending
	VDEV CUE EQU X'08'	Virtual control unit end
	VDEVNRDY EQU X'04'	Virtual device not ready
	VDEV CATT EQU X'02'	Virtual device attached by console function
	VDEVDED EQU X'01'	VDEVREAL is dedicated device RDEVBLOK
7 7	VDEVFLAG DS 1X	V*4 - virtual device flags
	Bits Defined in VDEVFLAG	
	VDEVRDO EQU X'80'	DASD - read-only
	VDEVENAB EQU X'80'	Virtual 270X - line enabled
	VDEV TDSK EQU X'40'	DASD - TDISK space allocated by CP
	VDEV DIAL EQU X'40'	VIRTUAL 270x - line connected
	VDEVCSPL EQU X'40'	Console - activity spooled
	VDEV231T EQU X'20'	DASD - 2311 simulated on top half of 2314
	VDEV231B EQU X'10'	DASD 2311 simulated on bottom half of 2314
	VDEVCCW1 EQU X'10'	Console and spooling - processing first CCW
	VDEV SAS EQU X'08'	DASD - Executing stand-alone seek
	VDEVRSRL EQU X'02'	Reserve/Release are valid CCW operation codes
	VDEVUC EQU X'01'	Virtual device sense bytes present
8 8	VDEVCSW DS 1D	Virtual channel status word

10	16	VDEVRELN	DS	1H	Virtual DASD cylinder relocation
12	18	VDEVBNB	DS	1H	Virtual DASD size (in cylinders)
14	20	VDEVPOSN	DS	1F	Virtual DASD seek position
18	24	VDEVQUED	DS	1F	Virtual SIO to real SIO queued time
1C	28	VDEVOPER	DS	1F	Device operational time
20	32	VDEVLINK	DS	1F	Link to virtual shared devices
24	36	VDEVREAL	DS	1F	Pointer to real device RDEVBLK
28	40	VDEVIOCT	DS	1F	Virtual Device I/O count
2C	44	VDEVUSER	DS	1F	Pointer to VMVLOK of VDEVBLK owner
30	48	VDEVIOER	DS	1F	Pointer to IOERBLK for last error
34	52	VDEVI0B	DS	1F	Pointer to active IOBLK
		VDEVSIZE	EQU	(*-VDEVBLK)/8	VDEVBLK size in doublewords (X'07')
					For spooling/console devices
		ORG		VDEVRELN	
10	16	VDEVXUSR	DS	CL8	Transferred to VMUSER
18	24	VDEVCON	DS	1F	Pointer to VCONCTL console control
1C	28	VDEVSPPL	DS	1F	Pointer to VSPLCTL spool control
20	32	VDEVCLAS	DS	1C	Spool - output class
21	33	VDEVKEY	DS	1X	Storage key in user's CAW
22	34	VDEVUNIT	DS	1H	Spool - output directed device address
24	36	VDEVCOPY	DS	1H	Number of copies requested
26	38	VDEVCLG	DS	1X	Console - virtual console flags
					Bits Defined in VDEVCLG
		VDEVATTN	EQU	X'80'	User pressed Attention more than once
		VDEVTIC	EQU	X'40'	Last CCW processed was a TIC
		VDEVTRAN	EQU	X'20'	Data transfer occurred during this channel program
		VDEVVCF	EQU	X'10'	Virtual console function in progress
		VDEVAUCR	EQU	X'08'	Auto carriage return on first read
27	39	VDEVSFGL	DS	1X	Spool - virtual spool flags
					Bits Defined in VDEVSFGL
		VDEVFEED	EQU	X'80'	Spool reader - last command was a feed
		VDEVXFER	EQU	X'80'	Spool output - transferred to VDEVXUSR
		VDEVCONT	EQU	X'40'	Spool input - continuous reading
		VDEVCP	EQU	X'40'	Spool output - continuous printing
		VDEVHOLD	EQU	X'20'	Hold output - save input
		VDEVEOF	EQU	X'08'	Spool input - set unit exception at EOF
		VDEVTERM	EQU	X'08'	Terminal output required for spooled console
		VDEVCFCL	EQU	X'04'	Device closed by console function
		VDEVPURG	EQU	X'02'	Spool output - purge file at close
		VDEVDIAG	EQU	X'02'	Spool input - device opened by DIAGNOSE
		VDEVSVCS	EQU	X'01'	Spool output - DMKVSP entered via SVC
		ORG		VDEVIOER	
30	48	VDEVSNSE	DS	1F	Sense bytes for spool device
34	52	VDEVFCBK	DS	1F	Address of forms control blok (VFCBBLOK)
		ORG		VDEVLINK	
20	32	VDEVTMAT	DS	1F	TDISK attached time (TOD clock word 0)

VFCBBLOK - VIRTUAL FORM CONTROL BUFFER BLOCK



Displacement			Field		Field Description, Contents, Meaning
Hex	Dec		Name		
0	0		VFCBCNT DS	1H	Current pointer to carriage column
2	2		VFCBFLAG DS	1X	V*1 working flag byte
			Bits Defined in VFCBFLAG		
			VFCBEOF EQU	X'80'	End of forms passed once
			VFCBCMD EQU	X'40'	forms control given
3	3		VFCBCHL DS	1X	V*2 Channel number or space count
4	4		VFCBWORK DS	1F	Work area
8	8		VFCBSPAR DS	2X	Spare
A	10		VCFBNDEX DS	1X	V*3 Index byte value
B	11		VFCBLOAD DS	CL181	Form control buffer area
			VFCBSIZE EQU	(*-VFCBBLOK)/8	Size in doublewords (X'18')

VMBLOK - VIRTUAL MACHINE CONTROL BLOCK

0	VMQFPNT		VMQBPNT					
8	VMPNT		VMECEXT					
10	VMSEG		VMSIZE					
18	VMCHSTRT		VMCUSTRT					
20	VMDVSTRT		VMTERM					
28	VMCHCNT	VMCUCNT	VMDVCNT	VMIOACTV				
30	VMCHTBL							
50	V*3	V*4	V*5	V*6	V*7	V*8	V*9	V*10
58	V*11	V*12	V*13	V*14	VMEXTINT	VMIOINT		
60	VMSLOCK	VMLLOCK		VMTIMER				
68	VMVTIME							
70	VMTMOUTQ							
78	VMTTIME							
80	VMTMINQ							
88	VMTODINQ							
90	VMINST			VMUPRIOR				
98	VMTREXT		VMADSTOP					

A0	VMPSW			
A8	VMGPRS			
E8	VMFPRS			
108	VMUSER			
110	VMACNT			
118	VMDIST			
120	VMPGREAD		VMPGWRIT	
128	VMWCNT	VMSEGDSP	VMSHRSYS	
130	VMIOCNT		VMPNCH	
138	VMLINS		VMCRDS	
140	VMCOMND			
148	VMTIMEON	VMDEDCH	VMQPRIOR	
150	VMPSWDCT	VMVTRMAD	VMPAGES	VMWSPROJ
158	VMTRQBLK	VMPRGIL	VMSTEALS	
160	VMPDRUM	VMPDISK		VMACOUNT
168	VMRDINQ		VMPPGRINQ	
170	VMEPRIOR		VMRSV4	
178	VMRSV5		VMRSV6	
180	VMUSER1		VMUSER2	
188	VMUSER3		VMUSER4	

Displacement  
Hex Dec

Field  
Name

0	0	VMQFPNT	DS	1F
4	4	VMQBPNT	DS	1F
8	8	VMPNT	DS	1F

Field Description, Contents, Meaning

Pointer to next VMBLOK in queue  
Pointer to previous VMBLOK in queue  
Pointer (CYCLIC) to next VMBLOK

C	12	VMCEXT	DS	1F	VMBLOK extended control pointer - ECBLOK
C	12	VMVCRO	EQU	VMCEXT	Virtual control register 0 for non-EC mode machine
10	16	VMSEG	DS	1F	Pointer to VMSEGTBL
14	20	VMSIZE	DS	1F	Virtual storage size - bytes
18	24	VMCHSTRT	DS	1F	Pointer to VCHBLOK table
1C	28	VMCUSTRT	DS	1F	Pointer to VCUBLOK table
20	32	VMDVSTRT	DS	1F	Pointer to VDEVBLOK table
24	36	VMTERM	DS	1F	Pointer to RDEVBLOK for user terminal
28	40	VMHCNT	DS	1H	Virtual channel count
2A	42	VMCUCNT	DS	1H	Virtual control unit count
2C	44	VMDVCNT	DS	1H	Virtual device count
2E	46	VMIOACTV	DS	1H	Active channel mask
30	48	VMCHTBL	DS	16H	Channels attached - VMCHSTRT index
50	80	VMRSTAT	DS	1X	V*3 - virtual machine running status
		Bits Defined in VMRSTAT			
		VMCFWAIT	EQU	X'80'	Waiting - Executing console function
		VMPGWAIT	EQU	X'40'	Waiting - paging operation(s)
		VMIOWAIT	EQU	X'20'	Waiting - Scheduled IOBLOK start
		VMPSWAIT	EQU	X'10'	Waiting - virtual PSW wait state
		VMEXWAIT	EQU	X'08'	Waiting - Instruction simulation
		VMLOGON	EQU	X'04'	User not logged on
		VMLOGOFF	EQU	X'02'	User logging off
		VMIDLE	EQU	X'01'	Virtual machine in idle wait state
		VMCPWAIT	EQU	VMCFWAIT+VMPGWAIT+VMIOWAIT+VMEXWAIT+VMLOGOFF+VMLOGON	
		VMNORUN	EQU	VMCPWAIT+VMPSWAIT	
		VMLONGWT	EQU	VMCFWAIT+VMLOGON+VMLOGOFF+VMIDLE	
51	81	VMDSTAT	DS	1X	V*4 - virtual machine dispatching status
		Bits Defined in VMDSTAT			
		VMDSP	EQU	X'80'	Virtual machine is dispatched runuser
		VMTSEND	EQU	X'40'	Virtual machine is compute bound
		VMTIO	EQU	X'10'	Virtual Machine is in TIO/SIO busy loop
		VMRUN	EQU	X'08'	Virtual machine runnable
		VMINQ	EQU	X'04'	Virtual machine in a queue
		VMELIG	EQU	X'02'	Reserved for IBM use
52	82	VMOSTAT	DS	1X	V*5 - virtual machine operating status
		Bits Defined in VMOSTAT			
		VMSYSOP	EQU	X'80'	Virtual machine is system operator
		VMSHR	EQU	X'40'	Virtual machine running shared system
		VMDISC	EQU	X'10'	Virtual machine console disconnected
		VMCFRUN	EQU	X'08'	Virtual machine running in CF mode
		VMVIRCF	EQU	X'04'	Virtual machine executing virtual CF
		VMCF	EQU	X'02'	Virtual machine executing CF
		VMKILL	EQU	X'01'	Virtual machine is to be logged off
53	83	VMQSTAT	DS	1X	V*6 - Virtual machine queueing status
		Bits Defined in VMQSTAT			
		VMPRIDSP	EQU	X'80'	Virtual eligible for Queue1



54 84

VMPSTAT DS 1X  
Bits Defined in VMPSTAT  
VMISAM EQU X'80'  
VMV370R EQU X'40'  
VMRPAGE EQU X'20'  
VMREAL EQU X'10'  
VMNOTRAN EQU X'08'  
VMPNMCS EQU X'04'  
VMACCOUN EQU X'02'

V\*7 - virtual machine processing status

Virtual machine has ISAM CCW checking  
Virtual machine can use extended format  
Virtual machine can reserve pages  
Virtual machine has V=R option  
No CCW translation for V=R user  
Reserved for IBM use  
Virtual machine may punch account cards



55	85	VMESTAT DS 1X Bits Defined in VMESTAT	V*8 - Virtual machine control status
		VMSHADI EQU X'80'	Shadow tables are present
		VMPERCM EQU X'40'	Virtual/CP PER active
		VMBADCRO EQU X'20'	Virtual control register 0 is invalid
		VMBADCR1 EQU X'10'	Virtual control register 1 is invalid
		VMEXTCM EQU X'08'	Virtual machine in extended control mode
		VMNEWCRO EQU X'04'	Virtual control register 0 has changed
		VMINVSEG EQU X'02'	All shadow tables invalid
		VMINVPAG EQU X'01'	Shadow page tables invalid
56	86	VMTRCTL DS 1X Bits Defined in VMTRCTL	V*9 - virtual machine tracing control
		VMTRPER EQU X'80'	Virtual PER tracing active
		VMTR SVC EQU X'40'	TRACE user SVC instructions
		VMTRPRG EQU X'20'	TRACE virtual program interrupts
		VMTRIO EQU X'10'	TRACE virtual I/O interrupts
		VMTREX EQU X'08'	TRACE external interrupts
		VMTRPRV EQU X'04'	TRACE user privileged instructions
		VMTRSIO EQU X'02'	TRACE virtual I/O instructions
		VMTRBRIN EQU X'01'	Trace successful branches or all instructions
		VMTRINT EQU VMTR SVC+VMTRPRG+VMTRIO+VMTREX	Trace all user interrupts
57	87	VMMLEVEL DS 1X Bits Defined in VMMLEVEL	V*10- message level
		VMMSGON EQU X'80'	Receiving messages
		VMWNGON EQU X'40'	Receiving warnings
		VMMCODE EQU X'20'	Receiving error message codes
		VMMTEXT EQU X'10'	Receiving texts of error messages
		VMMLINED EQU X'08'	Line editing on
		VMMACCCN EQU X'04'	Receiving accounting information
58	88	VMQLEVEL DS 1X Bits Defined in VMQLEVEL	V*11- queue level
		VMQ1 EQU X'80'	Virtual machine is interactive
		VMCOMP EQU X'40'	Virtual machine is compute bound
		VMHIPRI EQU X'20'	Virtual machine is highest priority
		VMLOPRI EQU X'10'	Virtual machine is lowest priority
		VMAEX EQU X'08'	Virtual machine is assured execution
		VMAEXP EQU X'04'	Virtual machine is assured percentage
		VMDROP1 EQU X'02'	Virtual machine just dropped from Q1
59	89	VMCLEVEL DS 1X Bits Defined in VMCLEVEL	V*12- command level
		VMCLASSA EQU X'80'	Class A functions
		VMCLASSB EQU X'40'	Class B functions
		VMCLASSC EQU X'20'	Class C functions
		VMCLASSD EQU X'10'	Class D functions
		VMCLASSE EQU X'08'	Class E functions

		VMCLASSF EQU	X'04'	Class F functions
		VMCLASSG EQU	X'02'	Class G functions
		VMCLASSH EQU	X'01'	Class H functions
5A	90	VMTLEVEL DS	1X	V*13- timer level
		Bits Defined in VMTLEVEL		
		VMTON EQU	X'80'	Virtual timer running
		VMRON EQU	X'40'	Virtual real timer running
		VMSTMPI EQU	X'08'	Virtual interval timer request queued
		VMSTMPT EQU	X'04'	Virtual CPU timer request queued
5B	91	VMPEND DS	1X	V*14- Interrupt pending summary flag
		Bits Defined in VMPEND		
		VMPERPND EQU	X'40'	Virtual PER interrupt pending
		VMPRGPND EQU	X'20'	Virtual program interrupt deferred
		VMSVCPND EQU	X'10'	Virtual SVC interrupt deferred
		VMICPND EQU	X'02'	Virtual I/O interrupt pending
		VMEXTPND EQU	X'01'	Virtual external interrupt pending
5C	92	VMEXTINT DS	1H	External interrupt pending flags
		Bits Defined in VMEXTINT		
		VMCKCINT EQU	X'08'	Clock comparator interrupt pending
		VMCPTINT EQU	X'04'	CPU timer interrupt pending
		Bits Defined in VMEXTINT+1		
		VMINTINT EQU	X'80'	Interval timer interrupt pending
		VMKEYINT EQU	X'40'	RED button interrupt pending
		VMSIGINT EQU	X'2F'	External signals pending
5E	93	VMIOINT DS	1H	I/O interrupt pending flags
60	96	VMSLOCK DS	1H	Short lock - reserved for IBM use
62	98	VMLLOCK DS	1H	Long lock - reserved for IBM use
64	100	VMTIMER DS	1F	Virtual timer value - X'50'
68	104	VMVTIME DS	1D	Virtual CPU time used - 2s complement
70	112	VMTHOUTQ DS	1D	VMTIME for exit from queue -2s complement
78	120	VMTTIME DS	1D	Total CPU time used - 2s complement
80	128	VMTMINQ DS	1D	VMTIME value at entry to queue
88	136	VMTODINQ DS	1D	TOD clock time stamp at queue entry
90	144	VMINST DS	XL6	Virtual machine privileged or tracing instruction
96	150	VMUPRIOR DS	1H	User priority from directory
98	152	VMTREXT DS	1F	Address of extended trace control block
9C	156	VMADSTOP DS	1F	Address of address stop control block
A0	160	VMPSW DS	1D	Virtual machine PSW
A8	168	VMGPRS DS	16F	Virtual machine general registers
E8	232	VMFPRS DS	4D	Virtual machine floating point registers
108	264	VMUSER DS	CL8	Virtual machine identification
110	272	VMACNT DS	CL8	Virtual machine accounting number
118	280	VMDIST DS	CL8	Virtual machine distribution code
120	288	VMPGREAD DS	1F	Total page reads
124	292	VMPGWRT DS	1F	Total page writes
128	296	VMWCNT DS	1H	Page wait count

12A	298	VMSEGDSP	DS	1H	Displacement of virtual machine SEGTABLE from start of block
12C	300	VMSHRSYS	DS	1F	Pointer to shared system table
130	304	VMIOCNT	DS	1F	Virtual SIO count for non-spooled I/O
134	308	VMPNCH	DS	1F	Virtual card count - spooled punch
138	312	VMLINS	DS	1F	Virtual line count - spooled printer
13C	316	VMCARDS	DS	1F	Virtual card count - spooled reader
140	320	VMCOMND	DS	CL8	Last CP command executed
148	328	VMTIMEON	DS	1F	LOGON time -TOD clock word 0
14C	332	VMDEDCH	DS	1H	Dedicated channel mask
14E	334	VMQPRIOR	DS	1H	Priority in dispatching queue
150	336	VMPSWDCT	DS	1H	Count of incorrect passwords entered
152	338	VMVTRMAD	DS	1H	Virtual terminal device address
154	340	VMPAGES	DS	1H	Number of pages currently resident
156	342	VMWSPROJ	DS	1H	Projected working set size
158	344	VMTRQBLK	DS	1F	Address of TRQBLOK for real timer
15C	348	VMPRGIL	DS	1H	ILC for pending program interrupt
15E	350	VMSTEALS	DS	1H	Number of waits for stolen pages
160	352	VMPDRUM	DS	1H	Reserved for IBM use
162	354	VMPDISK	DS	1H	Reserved for IBM use
164	356	VMACOUNT	DS	1F	Address of user ACCTBLOK
168	360	VMRDINQ	DS	1F	Page read total (VMPGREAD) at Q entry
16C	364	VMPGRINQ	DS	1F	Sum of VMPAGES count at each page read
170	368	VMEPRIOR	DS	1F	Eligible list priority
174	372	VMRSV4	DS	1F	Reserved for IBM use
178	376	VMRSV5	DS	1F	Reserved for IBM use
17C	380	VMRSV6	DS	1F	Reserved for IBM use
180	384	VMUSER1	DS	1F	Reserved for installation use
184	388	VMUSER2	DS	1F	Reserved for installation use
188	392	VMUSER3	DS	1F	Reserved for installation use
18C	396	VMUSER4	DS	1F	Reserved for installation use
		VMBSIZE	EQU	(*-VMBLOK)/8	VMBLOK size in doublewords (X'32')

VSPCTL - VIRTUAL SPOOL CONTROL BLOCK

0	VSPCAW		VSPDPAGE
8	VSPVPAGE		VSPRECNO
10	VSPNEXT	VSPIDACT	VSPSFBLK
18	VSPCCW		
20	VSPBUFBK		VSPMISC
28	V*1	VSPIDAL	VSPIDAW2

Displacement		Field		Field Description, Contents, Meaning
Hex	Dec	Name		
0	0	VSPCAW	DS 1F	Virtual address of user CCW
4	4	VSPDPAGE	DS 1F	DASD location (DCHR) of current page buffer
8	8	VSPVPAGE	DS 1F	Virtual address of page buffer
C	12	VSPRECNO	DS 1F	Records remaining in current buffer
10	16	VSPNEXT	DS 1H	DISP. in buffer of next record start
12	18	VSPIDACT	DS 1H	Data byte count of IDA CCW
14	20	VSPSFBLK	DS 1F	Pointer to SFBLOK for file
18	24	VSPCCW	DS 1D	Current user CCW
20	32	VSPBUFBK	DS 1F	Address of a buffer area
24	36	VSPMISC	DS 1F	Use varies according to caller
28	40	VSPIDASW	DS 1X	V*1 IDA work flag
29	41	VSPIDAL	DS 3X	Address of indirect data list
2C	44	VSPIDAW2	DS 1F	Contains IDAW2
		VSPSIZE	EQU (*-VSPCTL)/8	Size in doublewords (X'06')
		VSPBUFSZ	EQU (200)/8	Size in doublewords (X'19')

DIAGNOSTIC AIDS

COMMAND-TO-MODULE CROSS-REFERENCE

<u>Command</u>	<u>Entry Point</u>	<u>Messages</u>			
					DMKVDB142E DMKVDB143E
			BACKSPACE	DMKCSOBS	DMKCS0003E DMKCS0006E DMKCS0021E DMKCS0040E DMKCS0046E DMKCS0140E DMKCS0141E
ACNT	DMKCPVAC	DMKCPV003E DMKCPV007E DMKCPV020E DMKCPV045E			
ADSTOP	DMKCFDAD	DMKCFD004E DMKCFD026E DMKCFD160E DMKCFD161E	BEGIN  CHANGE	DMKCFMBE  DMKCSUCH	DMKCFM004E  DMKCSU003E DMKCSU006E DMKCSU008E DMKCSU013E DMKCSU026E DMKCSU027E DMKCSU028E DMKCSU029E DMKCSU030E DMKCSU032E DMKCSU035E DMKCSU042E
ATTACH (channel)	DMKVCHDC	DMKVCH034E DMKVCH048E DMKVCH129E DMKVCH131E DMKVCH132E			
	DMKVDBAT	DMKVDB020E DMKVDB045E			
ATTACH	DMKVDBAT	DMKVDB003E DMKVDB006E DMKVDB020E DMKVDB021E DMKVDB022E DMKVDB023E DMKVDB034E DMKVDB040E DMKVDB045E DMKVDB046E DMKVDB120E DMKVDB122E DMKVDB123E DMKVDB124E DMKVDB125E DMKVDB126E DMKVDB127E DMKVDB128E DMKVDB133E DMKVDB134E	CLOSE  COUPLE  DCP	DMKCSPCL  DMKDIACP  DMKCDBDC	DMKCSP003E DMKCSP006E DMKCSP013E DMKCSP022E DMKCSP029E DMKCSP032E DMKCSP040E  DMKDIA006E DMKDIA011E DMKDIA020E DMKDIA022E DMKDIA040E DMKDIA045E DMKDIA047E DMKDIA058E  DMKCDB003E

		DMKCDB004E			DMKCPV046E
		DMKCDB009E			DMKCPV140E
		DMKCDB010E			
		DMKCDB026E	DISCONN	DMKUSODS	DMKUSO003E
		DMKCDB033E			
		DMKCDB160E	DISPLAY	DMKCDBDI	DMKCDB003E
DEFINE	DMKDEFIN	DMKDEF003E			DMKCDB004E
		DMKDEF022E			DMKCDB009E
		DMKDEF024E			DMKCDB010E
		DMKDEF025E			DMKCDB026E
		DMKDEF026E			DMKCDB160E
		DMKDEF040E	DMCP	DMKCDBDM	DMKCDB003E
		DMKDEF091E			DMKCDB004E
		DMKDEF092E			DMKCDB009E
		DMKDEF094E			DMKCDB033E
		DMKDEF136E			DMKCDB160E
DETACH (channel)	DMKVCHDC	DMKVCH034E	DRAIN	DMKCSODR	DMKCSO003E
		DMKVCH048E			DMKCSO006E
		DMKVCH130E			DMKCSO021E
	DMKVDBDE	DMKVDB020E			DMKCSO040E
		DMKVDB034E			DMKCSO046E
					DMKCSO140E
DETACH	DMKVDBDE	DMKVDB006E	DUMP	DMKCDBDU	DMKCDB003E
		DMKVDB020E			DMKCDB004E
		DMKVDB021E			DMKCDB009E
		DMKVDB022E			DMKCDB033E
		DMKVDB040E			DMKCDB160E
		DMKVDB045E			
		DMKVDB046E			
		DMKVDB121E	ECHO	DMKMSGEC	none.
		DMKVDB123E			
		DMKVDB124E	ENABLE	DMKCPVEN	DMKCPV003E
		DMKVDB135E			DMKCPV006E
		DMKVDB140E			DMKCPV021E
DIAL	DMKDIAL	DMKDIA011E			DMKCPV026E
		DMKDIA020E			DMKCPV040E
		DMKDIA022E			DMKCPV046E
		DMKDIA045E			DMKCPV140E
		DMKDIA047E	EXTERNAL	DMKCPBEX	DMKCPB005E
		DMKDIA055E			
		DMKDIA056E	FLUSH	DMKCSOFL	DMKCSO003E
DISABLE	DMKCPVDS	DMKCPV003E			DMKCSO006E
		DMKCPV006E			DMKCSO013E
		DMKCPV021E			DMKCSO021E
		DMKCPV026E			DMKCSO040E
		DMKCPV040E			DMKCSO046E
					DMKCSO140E



		DMKCS0141E			DMKLNK107E
FORCE	DMKUSOFL	DMKUSO003E DMKUSO020E DMKUSO045E			DMKLNK108E DMKLNK109E DMKLNK110E DMKLNK111E DMKLNK112E DMKLNK113E DMKLNK114E DMKLNK115E DMKLNK116E DMKLNK117E DMKLNK137E
FREE	DMKCSPPR	DMKCSP006E DMKCSP007E DMKCSP020E DMKCSP053E			
HALT	DMKCPVH	DMKCPV021E DMKCPV040E DMKCPV144W	LOADBUF	DMKCSOLD	DMKCS0003E DMKCS0006E DMKCS0013E DMKCS0021E DMKCS0026E DMKCS0031E DMKCS0036E DMKCS0040E DMKCS0043E DMKCS0046E DMKCS0140E DMKCS0142E DMKCS0148E
HOLD	DMKCSPHL	DMKCSP006E DMKCSP007E DMKCSP020E DMKCSP053E			
IPL	DMKCFPIP	DMKCFP002E DMKCFP003E DMKCFP013E DMKCFP022E DMKCFP026E DMKCFP040E DMKCFP044E DMKCFP170E DMKCFP171E DMKCFP172E DMKCFP173E DMKCFP174E DMKCFP177E DMKVMI230E DMKVMI231E DMKVMI232E DMKVMI233E DMKVMI234E	LOADVFCB	DMKCSOVL	DMKCS0006E DMKCS0022E DMKCS0026E DMKCS0031E DMKCS0036E DMKCS0040E DMKCS0043E
			LOCATE	DMKCFDLO	DMKCFD021E DMKCFD022E DMKCFD026E DMKCFD040E
LINK	DMKLNKIN	DMKLNK020E DMKLNK022E DMKLNK052E DMKLNK053E DMKLNK101W DMKLNK102W DMKLNK103W DMKLNK104E DMKLNK105E DMKLNK106E	LOCK	CMKCPVLK	DMKCPV004E DMKCPV009E DMKCPV020E DMKCPV033E DMKCPV045E DMKCPV160E
			LOGOFF (LOGOUT)	DMKUSOLG	DMKUSO003E

LOGON (LOGIN, operator)	DMKLOGOP	DMKLOG003E DMKLOG020E			DMKCQP020E DMKCQP021E
LOGON (LOGIN, user)	DMKLOGON	DMKLOG050E DMKLOG051E DMKLOG052E DMKLOG053E DMKLOG054E DMKLOG090E DMKLOG091E DMKLOG092E DMKLOG093E	READY	DMKCPBRY	DMKCQP022E DMKCQP040E DMKCQP045E DMKCPB006E DMKCPB022E DMKCPB040E
			REPEAT	DMKCSORP	DMKCS0003E DMKCS0006E DMKCS0013E DMKCS0021E DMKCS0030E DMKCS0040E DMKCS0046E DMKCS0140E DMKCS0141E
MONITOR	DMKMCCCL	DMKMCC002E DMKMCC026E			
MSG	DMKMSGMS	DMKMSG003E DMKMSG020E DMKMSG045E DMKMSG057W			
NOTREADY	DMKCPBNR	DMKCPB006E DMKCPB022E DMKCPB040E	RESET	DMKCPBRS	DMKCPB022E DMKCPB040E
			REWIND	DMKCPBRW	DMKCPB006E DMKCPB022E DMKCPB040E DMKCPB059E
ORDER	DMKCSUOR	DMKCSU003E DMKCSU006E DMKCSU008E DMKCSU026E DMKCSU027E DMKCSU028E DMKCSU035E DMKCSU042E		DMKCFGSV	DMKCFG026E DMKCFG044E DMKCFG170E DMKCFG171E DMKCFG172E DMKCFG173E DMKCFG435E
			SAVESYS		
PURGE	DMKCSUPU	DMKCSU003E DMKCSU006E DMKCSU008E DMKCSU026E DMKCSU028E DMKCSU035E DMKCSU042E	SET	DMKCFSET	DMKCFPS003E DMKCFPS006E DMKCFPS013E DMKCFPS021E DMKCFPS026E DMKCFPS040E DMKCFPS041E DMKCFPS045E DMKCFPS046E DMKCFPS140E DMKCFPS175E
QUERY (initialize) (Class G)	DMKCFMQU DMKCQGEN	DMKCFM026E DMKCQG020E DMKCQG022E DMKCQG027E DMKCQG040E DMKCQG042E DMKCQG045E			DMKMCH003E DMKMCH026E
(Class B, E, G)	DMKCQPRV	DMKCQP003E DMKCQP006E		DMKMCHMS	

SHUTDOWN	DMKCPVSH	none.	SYSTEM	DMKCPBSR	DMKCPB012E DMKCPB026E
SLEEP	DMKCFMSL	none.	TERMINAL	DMKCFTRM	DMKCFT002E DMKCFT006E DMKCFT026E
SPACE	DMKCSOSP	DMKCSO006E DMKCSO021E DMKCSO040E DMKCSO046E DMKCSO140E DMKCSO141E	TRACE	DMKTRA	DMKTRA002E DMKTRA003E DMKTRA013E DMKTRA026E DMKTRA180E DMKTRA181E
SPOOL	DMKCSPSP	DMKCSP003E DMKCSP006E DMKCSP007E DMKCSP013E DMKCSP020E DMKCSP022E DMKCSP026E DMKCSP028E DMKCSP030E DMKCSP040E DMKCSP053E	TRANSFER	DMKCSUTR	DMKCSU003E DMKCSU007E DMKCSU008E DMKCSU020E DMKCSU026E DMKCSU027E DMKCSU028E DMKCSU042E DMKCSU053E
START	DMKCSOSD	DMKCSO003E DMKCSO006E DMKCSO013E DMKCSO021E DMKCSO028E DMKCSO040E DMKCSO046E DMKCSO140E	UNLOCK	DMKCPVUL	DMKCPV004E DMKCPV009E DMKCPV020E DMKCPV033E DMKCPV045E DMKCPV160E DMKCPV176E DMKCPV202E
STCP	DMKCDSCP	DMKCDS004E DMKCDS005E DMKCDS026E DMKCDS033E DMKCDS160E DMKCDS162E	VARY	DMKCPVRY	DMKCPV003E DMKCPV021E DMKCPV026E DMKCPV040E DMKCPV049E DMKCPV123E DMKCPV124E DMKCPV140E DMKCPV142E
STORE	DMKCDSTO	DMKCDS004E DMKCDS005E DMKCDS010E DMKCDS012E DMKCDS026E DMKCDS033E DMKCDS160E DMKCDS161E DMKCDS162E DMKCDS163E	WNG	DMKMSGWN	DMKMSG003E DMKMSG020E DMKMSG045E DMKMSG057W

MESSAGE-TO-FLOWCHART CROSS-REFERENCE

DMKACO425A	4	DMKCFM001E	2
DMKBLD200E	1	DMKCFM004E	4
DMKBLD201E	1	DMKCFM026E	5
DMKBLD202E	1	DMKCFP002E	6
DMKCC601I	4	DMKCFP003E	5
DMKCC602I	4	DMKCFP013E	6
DMKCC603W	1	DMKCFP026E	9
DMKCC604I	3	DMKCFP040E	7
DMKCC605I	2	DMKCFP044E	5
DMKCC606I	4 (in module DMKIOE)	DMKCFP170E	8
DMKCDB003E	4	DMKCFP171E	8
DMKCDB004E	7	DMKCFP172E	8
DMKCDB009E	8	DMKCFP173E	8
DMKCDB010E	7	DMKCFP174E	7
DMKCDB026E	1	DMKCFP177E	10
DMKCDB033E	3	DMKCFM003E	1
DMKCDB160E	8	DMKCFM006E	9
DMKCDSD004E	3	DMKCFM013E	13
DMKCDSD005E	3	DMKCFM021E	7
DMKCDSD010E	5	DMKCFM026E	1
DMKCDSD012E	4	DMKCFM040E	7
DMKCDSD026E	1	DMKCFM041E	10
DMKCDSD033E	3	DMKCFM045E	4
DMKCDSD160E	2	DMKCFM046E	9
DMKCDSD161E	2	DMKCFM140E	9
DMKCDSD162E	5	DMKCFM175E	5
DMKCDSD162W	6	DMKCFM002E	1
DMKCDSD163E	5	DMKCFM006E	5
DMKCFD004E	3	DMKCFM026E	1
DMKCFD021E	2	DMKCKP900E	7
DMKCFD022E	1	DMKCKP901E	7
DMKCFD026E	1	DMKCKP902E	3
DMKCFD040E	1	DMKCKP904E	5
DMKCFD160E	3	DMKCKP910I	3
DMKCFD161E	3	DMKCKP911W	3
DMKCFG026E	1	DMKCKP912W	3
DMKCFG044E	1	DMKCKP960I	8
DMKCFG170E	1	DMKCKP961W	8
DMKCFG171E	2	DMKCNS454I	7
DMKCFG172E	2	DMKCNS455I	1
DMKCFG173E	2	DMKCNS500I	21
DMKCFG435E	3	DMKCNS501I	21
		DMKCNS502I	21
		DMKCNS503I	21

DMKCNS504I	21
DMKCNS505I	21
DMKCNS527I	21
DMKCNS528I	22
DMKCPB005E	2
DMKCPB006E	3
DMKCPB012E	2
DMKCPB022E	4
DMKCPB026E	1
DMKCPB040E	4
DMKCPB059E	4
DMKCPI950A	4
DMKCPI951I	4
DMKCPI952I	4
DMKCPI953I	5
DMKCPI954E	1
DMKCPI955W	2
DMKCPI960I	6
DMKCPI961W	6
DMKCPV003E	7
DMKCPV004E	4
DMKCPV006E	2
DMKCPV007E	6
DMKCPV009E	5
DMKCPV020E	4
DMKCPV021E	1
DMKCPV026E	1
DMKCPV033E	4
DMKCPV040E	1
DMKCPV045E	4
DMKCPV046E	2
DMKCPV049E	8
DMKCPV123E	9
DMKCPV124E	8
DMKCPV140E	2, 8
DMKCPV142E	8
DMKCPV144W	10
DMKCPV160E	5
DMKCPV176E	4
DMKCPV202E	4
DMKCQG020E	7
DMKCQG022E	6
DMKCQG027E	3
DMKCQG040E	7
DMKCQG042E	4
DMKCQG045E	7

DMKCQP003E	2
DMKCQP006E	4
DMKCQP020E	3
DMKCQP021E	4
DMKCQP022E	4
DMKCQP040E	4
DMKCQP045E	3
DMKCSO003E	1
DMKCSO006E	1
DMKCSO013E	1
DMKCSO021E	12
DMKCSO022E	12
DMKCSO026E	13
DMKCSO028E	3
DMKCSO030E	6
DMKCSO031E	8
DMKCSO036E	10
DMKCSO040E	12
DMKCSO043E	8
DMKCSO046E	1
DMKCSO140E	1
DMKCSO141E	2
DMKCSO142E	7
DMKCSO148E	11
DMKCSP003E	1
DMKCSP006E	1
DMKCSP007E	5
DMKCSP013E	1
DMKCSP020E	5
DMKCSP022E	7
DMKCSP026E	4
DMKCSP028E	8
DMKCSP029E	10
DMKCSP030E	9
DMKCSP032E	12
DMKCSP040E	8
DMKCSP053E	5
DMKCSU003E	4
DMKCSU006E	1
DMKCSU007E	10
DMKCSU008E	15
DMKCSU013E	2
DMKCSU020E	8
DMKCSU026E	4
DMKCSU027E	1
DMKCSU028E	12
DMKCSU029E	14

DMKCSU030E 13  
 DMKCSU032E 3  
 DMKCSU035E 12  
 DMKCSU042E 4  
 DMKCSU053E 11

DMKDAS500I 4  
 DMKDAS501A 4  
 DMKDAS502D 4  
 DMKDAS503I 2  
 DMKDAS504D 4  
 DMKDAS505D 4  
 DMKDAS506I 5  
 DMKDAS507D 4  
 DMKDAS508I 7  
 DMKDAS509I 5  
 DMKDAS513I 5  
 DMKDAS514D 2,4  
 DMKDAS516I 5  
 DMKDAS517I 6  
 DMKDAS518I 6  
 DMKDAS520I 2  
 DMKDAS956A 3

DMKDEF003E 3  
 DMKDEF022E 1  
 DMKDEF024E 2  
 DMKDEF025E 1  
 DMKDEF026E 1  
 DMKDEF040E 4  
 DMKDEF091E 2  
 DMKDEF092E 3  
 DMKDEF094E 6  
 DMKDEF136E 2

DMKDIA006E 8  
 DMKDIA011E 3  
 DMKDIA020E 1  
 DMKDIA022E 1  
 DMKDIA040E 9  
 DMKDIA045E 1  
 DMKDIA047E 1  
 DMKDIA055E 2  
 DMKDIA056E 2  
 DMKDIA058E 9

DMKDMP905W 5  
 DMKDMP906W 5  
 DMKDMP907W 3  
 DMKDMP908I 1

DMKDMP909W 8  
 DMKDSP450W 4  
 DMKDSP451W 1  
 DMKDSP452W 3

DMKIOF550E 4  
 DMKIOF551E 5  
 DMKIOF552E 12  
 DMKIOF553E 5

DMKIOG554E 3  
 DMKIOG555E 3  
 DMKIOG556I 4  
 DMKIOG557I 4  
 DMKIOG558I 6  
 DMKIOG559I 5  
 DMKIOG560I 5

DMKLNK020E 1  
 DMKLNK022E 2  
 DMKLNK052E 4  
 DMKLNK053E 2  
 DMKLNK101W 10  
 DMKLNK102W 10  
 DMKLNK103W 10  
 DMKLNK104E 10  
 DMKLNK105E 11  
 DMKLNK106E 11  
 DMKLNK107E 3  
 DMKLNK108E 8  
 DMKLNK109E 4  
 DMKLNK110E 8  
 DMKLNK111E 8  
 DMKLNK112E 8  
 DMKLNK113E 8  
 DMKLNK114E 6  
 DMKLNK115E 1  
 DMKLNK116E 5  
 DMKLNK117E 5  
 DMKLNK137E 5

DMKLOG003E 2  
 DMKLOG020E 1  
 DMKLOG050E 2  
 DMKLOG051E 1  
 DMKLOG052E 7  
 DMKLOG053E 1  
 DMKLOG054E 9  
 DMKLOG090E 8

DMKLOG091E	6
DMKLOG092E	5
DMKLOG093E	6
DMKMCC002E	1
DMKMCC026E	1
DMKMCH003E	8
DMKMCH026E	8
DMKMCH610I	2
DMKMCH611I	1,7
DMKMCH612W	1
DMKMCH613I	3
DMKMCH614I	4
DMKMCH615I	7
DMKMCH616I	3
DMKMCH617I	7
DMKMCH618I	8
DMKMCH619I	3
DMKMID453I	1
DMKMSG003E	1
DMKMSG020E	1
DMKMSG045E	1
DMKMSG057W	2
DMKPAG415E	3
DMKPGT400I	4
DMKPGT401I	4
DMKPRG453W	3
DMKPTR410W	2
DMKRSE500I	7
DMKRSE501A	7
DMKRSE501I	7
DMKRSE502I	7
DMKRSE503A	7
DMKRSE503I	7
DMKRSE504A	7
DMKRSE504I	7
DMKRSE505A	7
DMKRSE508I	7
DMKRSE520A	7
DMKRSE520I	7
DMKRSE521I	7
DMKRSE524I	7

DMKRSE525I	7
DMKRSE529I	7
DMKRSP426E	1
DMKRSP428E	3
DMKRSP430A	13
DMKRSP431A	11
DMKRSP432A	11
DMKRSP433A	11
DMKRSP434A	9
DMKSAV350W	2
DMKSAV351W	2
DMKSAV352W	3
DMKTAP500I	4
DMKTAP501A	4
DMKTAP502D	2
DMKTAP503I	2
DMKTAP504D	5,7
DMKTAP505D	4
DMKTAP510I	6
DMKTAP511I	2
DMKTAP512I	5
DMKTAP513I	6
DMKTAP516I	7
DMKTAP517I	2
DMKTAP518I	8
DMKTAP519I	6
DMKTAP520I	2
DMKTAP521I	4
DMKTAP522I	5
DMKTAP523I	5
DMKTRA002E	1
DMKTRA003E	1
DMKTRA013E	2
DMKTRA026E	1
DMKTRA180W	4
DMKTRA181E	3
DMKTRA182E	1
DMKUDR475I	2
DMKUSO003E	5
DMKUSO020E	2
DMKUSO045E	2
DMKVCH034E	1
DMKVCH048E	1

DMKVCH129E	2
DMKVCH130E	5
DMKVCH131E	2
DMKVCH132E	2
DMKVDB003E	3
DMKVDB006E	6
DMKVDB020E	2
DMKVDB021E	1
DMKVDB022E	1
DMKVDB023E	3
DMKVDB034E	1
DMKVDB040E	9
DMKVDB045E	10
DMKVDB046E	10
DMKVDB120E	5
DMKVDB121E	8
DMKVDB122E	11
DMKVDB123E	8
DMKVDB124E	11
DMKVDB125E	6
DMKVDB126E	6
DMKVDB127E	11
DMKVDB128E	7
DMKVDB133E	13

DMKVDB134E	11
DMKVDB135E	11
DMKVDB140E	10
DMKVDB142E	8
DMKVDB143E	8
DMKVM1110E	1
DMKVM1230E	2
DMKVM1231E	7
DMKVM1232E	2
DMKVM1233E	3
DMKVM1234E	2
DMKVSP427I	6
DMKVSP429I	5
DMKWRM902E	4
DMKWRM903E	2
DMKWRM904E	3
DMKWRM909E	2
DMKWRM910I	2
DMKWRM911W	2
DMKWRM920I	4



CP WAIT STATE CODES

The wait state code is found in the right half of the program status word (PSW) when the CPU is in wait state. A wait state is produced by one of the following modules:

- DMKCCH
- DMKCKP
- DMKCPI
- DMKDMP
- DMKMCH
- DMKSAV
- DMKWRM

When a wait state occurs, the PSW is displayed at the operator's console in the following format:

|        xxxxxxxxxxxzzzzzwww

Where:

xxxxxxxx is the left half of the program status word.

This half may be either:

03yyyyyy Valid wait condition. The system is waiting for work.

00yyyyyy System wait caused by an error condition.

zzzzzwww is the right half of the program status word. The wait state code, www, indicates the error condition.

Wait

Codes    Explanation

001        The machine check handler has encountered an irrecoverable failure. Probable hardware error.

002        The channel check handler has encountered an irrecoverable failure. Probable hardware error.

- 003        A system failure has occurred before a valid warm start was performed.
- 004        This wait state code is loaded by DMKDMP when a console, or an output device is not operational, or when a console or output device produce an inexplicable error status. Probable hardware error.
- 005        DMKCPI could not find an operational primary or alternate console. Probable hardware error.
- 006        This is a normal wait when a system shutdown is completed.
- 007        A program check, a machine check, or a permanent I/O error was encountered by the checkpoint program.
- 008        Checkpoint and system shutdown are complete.
- 009        An error condition has occurred that prevents a warm start.
- 00A        A machine check occurred while DMKSAV was attempting to save or restore a page image copy of the nucleus on a SYSRES device. Probable hardware error.
- 00B        A machine check occurred before initiation was complete.
- 00C        An attempt was made to IPL from a disk that did not contain a system. Thus, the wait code 00C enter on disk by the Format program is encountered.
- 00D        The size defined during system generation is greater than the real machine size, or a hardware error has occurred which inhibits VM/370 from using the required storage.
- 00F        Hardware errors are being received on VM/370 paging device(s). This wait state is preceded by message DMKPAG415E - CONTINUOUS PAGING ERRORS FROM DASDxxx.

ABEND CODES

<u>Code</u>	<u>Reason</u>
ELD001	An invalid pointer to the RDEVBLK was found in register 8 when DMKBLDVM was called to build a new VMBLOK.
CFM001	No stacked CPEXBLOK was found for a user with a pending LOGOFF flag set.
CNS001	Condition code 2 was returned by a TIO instruction to a logged on user communication line or console.
CNS002	Condition code 2 or 3 was returned from a SIO instruction to a logged on user communication line or console.
CNS003	Condition code 1 was returned from a SIO instruction to a logged on user communication line or console, accompanied by CSW status other than Attention or Unit Check.
CNS004	The input data count, less idle and control characters, for a read from a 2741 is less than 0.
CNS005	The input data count for a read from a 2741 is equal to 0.
CNS006	The data count at entry to the code translation routine is less than or equal to 0.
CNS007	The input data count less idle and pad characters is less than zero for a non-IBM terminal.
CNS008	The IOBIOER field of the IOBLOK contains an invalid pointer to an IOERBLOK.
CPI001	The RDEVBLK for the DASD on which the SYSRES volume is mounted cannot be located. The SYSRES volume is specified in the SYSRES macro in the module DMKSYS.
CPI002	A valid system directory file could not be located.
CPI003	The system TOD clock is not operational.
CVT001	The system TOD clock is in error or is not operational.
DRD001	The device code index in the compressed DASD address for the system dump file points to a RDEVLBOK for an invalid DASD. The valid DASDs are 2305, 3330, or 2314/2319.
DSP001	During I/O Interrupt Unstack and Reflection, DMKSCNVU could not locate all of the virtual control blocks for the interrupting unit.
DSP002	The dispatcher (DMKDSP) is attempting to dispatch a virtual relocate user whose shadow segment tables or virtual extended control register 0 are invalid.
DSP003	The dispatcher has sensed that the interval timer did not decrement properly.
DSP004	A virtual device was detached while an I/O interrupt was being traced and its VDEVBLK cannot be found.
FRE001	The size of the block being returned (via register 0) is less than or equal to 0.
FRE002	The address of the free storage block being returned matches the address of a block in the free storage chain.
FRE003	The address of the free storage block being returned overlaps the next lower block on the free storage chain.
FRE004	The address of the free storage block being returned overlaps the next higher block on the free storage chain.
FRE005	A module is attempting to release storage in the resident CP nucleus.
FRE006	A module is requesting a block of storage whose size (in register 0) is less than or equal to zero.

FRE007	A module is attempting to release a block of storage whose address exceeds the size of real storage.	PGT004	The dummy RECBLOK indicating the spooling DASD pages on the cylinder that are to be released contains a page count greater than the number of pages allocated on the cylinder.
FRE008	The address of the free storage block being returned matches the address of the first block in the subpool for that size.	PGT005	A module is attempting to release a DASD page slot on a cylinder for which no page allocation block (RECBLOK) exists.
FRE009	The address of the free storage block being returned matches the address of the second block in the subpool for that size.	PGT006	The last DASD page slot in a RECBLOK has been deallocated but the bit representing the cylinder in the cylinder allocation block (ALOCBLOK) is not set to one, indicate that the cylinder was not allocated.
FRE010	A program is attempting to extend free storage while storage is in the process of being extended.	PGT007	A module is attempting to release a page of virtual storage being used by CP that has not been marked allocated.
FRE011	A CP module has attempted to return a block of storage that is in the user dynamic paging area.	PRG001	Program check (operation) in the control program.
HVC001	The user pointed to by register 11 issued a diagnose while attempting to format the I/O Error or Channel Check/Machine Check recording areas: the SYSRES device type is unrecognizable.	PRG002	Program check (privileged operation) in the control program.
IOS001	The caller is attempting to reset an active IOBLOK that contains an invalid unit address.	PRG003	Program check (execute) in the control program.
IOS002	DMKIOS is attempting to restart an IOBLOK from the RCHBLOK queue, but that IOBLOK contains an invalid unit address.	PRG004	Program check (protection) in the control program.
IOS003	DMKIOS is attempting to remove an IOBLOK from a queue, but that IOBLOK is on more than one queue.	PRG005	Program check (addressing) in the control program.
PGT001	The number of cylinders in use stored in the allocation block (ALOCBLOK) is less than the maximum but DMKPGT was unable to find available cylinders.	PRG006	Program check (specification) in the control program.
PGT002	The count of pages in use in a page allocation block (RECBLOK) is less than the maximum but DMKPGT was unable to find available pages.	PRG007	Program check (data) in the control program.
PGT003	The DASD page slot being released is not marked allocated.	PRG008	Program check (fixed-point overflow) in the control program.
		PRG009	Program check (fixed point divide) in the control program.
		PRG010	Program check (decimal overflow) in the control program.

PRG011	Program check (decimal divide) in the control program.	PTR007	No storage available for extend.
PRG012	Program check (exponential overflow) in the control program.	PTR008	The CORTABLE entry on the free list points to a page currently active.
PRG013	Program check (exponential underflow) in the control program.	PTR009	The adjusted count of resident shared pages has fallen below zero.
PRG014	Program check (significance) in the control program.	PTR010	The adjusted count of resident reserved pages has fallen below zero.
PRG015	Program check (floating-point divide) in the control program.	QCN001	A HIO instruction attempting to halt a prepare/read to a logged on user terminal received a condition code of 3.
PRG254	A translation specification exception has been received for a virtual machine that is not in Extended Control Mode.	QCN002	A CP routine has attempted to initiate a read or write to a device whose RDEVTYPE field is invalid.
PRG255	A PER interrupt has been received for a virtual machine that is running with PER disabled in its virtual PSW.	RPA001	The virtual address supplied to DMKRPAGT is outside of the virtual storage being referenced.
PSA001	Free storage is not available for the save areas.	RPA002	The virtual address supplied to DMKPRAPT is outside of the virtual storage being referenced.
PSA002	The System Restart key in the CPU console was depressed.	RPA003	User page count is negative.
PSA003	Fatal DASD I/O error on paging device.	SCH001	The adjusted count of users in the in-queue (interactive plus non-interactive) has fallen below zero.
PTR001	A segment exception or a translation specification exception has occurred while executing a LRA (Load Real Address) instruction in DMKPTR.	TDK001	A program is attempting to deallocate a cylinder of T-disk space for which no cylinder allocation block (ALOCBLOK) exists.
PTR002	A program is attempting to unlock a page frame whose address exceeds the size of real storage.	TDK002	A program is attempting to deallocate a cylinder or cylinders of T-disk space that are not marked allocated.
PTR003	A program is attempting to unlock a real storage page frame whose CORTABLE entry is not flagged as locked.	TRC001	An erroneous call to TRACE was detected.
PTR004	The lock count in the CORTABLE entry for the page frame being unlocked has been decremented to a value that is less than 0.	UDR001	The user directory module is looping, trying to read all of the UDIRBLOK page buffers from the directory device. Or, a directory containing over 10,816 users was loaded.
PTR006	Request to extend storage while extending.		

VAT001	A hardware page exception occurred, but the translation tables indicate either the segment is not available or the page table does not exist.	VI0002	DMKSCNVU was unable to locate all of the virtual I/O control blocks for the virtual unit address associated with the interrupt previously unstacked.
VDB001	The VDEVBLOK for the virtual device being released contains an unrecognizable device type.	VI0003	DMKIOS has returned an IOBLOK indicating a condition code of 2 was received from the Start I/O for the operation.
VDB002	The 'Sysownd' list is in an invalid format.		
VDB003	DASD link chain is invalid.	VSP001	DMKSCNVU was unable to locate all of the virtual I/O control blocks for the channel program that was previously executed. The virtual I/O configuration was destroyed.
VI0001	The VMINST field in the user's VMBLOK issuing privileged I/O operation does not contain a recognizable I/O operation code.		



## APPENDIX A: VM/370 MODULE FORMAT

Every module for VM/370 is formatted in the following manner:

MOD            TITLE Card  
              ISEQ 73, 80    Validate source seq.

### Module Prologue

The prologue contains a heading for each of the topics listed in the order that they appear below, even if the topic does not apply to the given module. Topic headings start in column three and are followed by one blank line. The text beneath each heading should start in column ten. The required topics are as follows:

MODULE NAME - The actual name of the module (that is, the label of the START or first CSECT card)

FUNCTION - A brief (one or two sentence) description of the purpose of the module

ATTRIBUTES - A list of things such as whether or not the module is reentrant or serially reusable, resident or pageable, and how it is called (via an SVC, EALR, or GOTO).

ENTRY POINTS - A list of the name of each entry point, followed by a short explanation of the reason that the entry point is called. Also included is a list of all fields that are referenced or modified by an external routine.

ENTRY CONDITIONS - State any registers that must be loaded by the caller and what values they contain. If general registers contain parameters, state their symbolic values and their meanings.

EXIT CONDITIONS, NORMAL AND ERROR - State any of the caller's registers that are modified and what values they contain, if a meaningful condition code is set,

and indicate if the module does not return to its caller.

CALLS TO OTHER ROUTINES - A list of any external routines or modules that are called. Include any exits to DISPATCH via a GOTO.

EXTERNAL REFERENCES - A list of any tables, control blocks or values that are referenced in this module but are defined elsewhere.

TABLES/WORKEAREAS - A list of any temporary work or scratch areas. For example, the use of the BALRSAVE area in PSA or the SAVEWRK areas in the standard SAVEAREA for any purpose other than normal register saving and restoring.

REGISTER USAGE - A list of the usage for each general register that has a consistent purpose through the module. Also list any unused registers and any registers used only for scratch or intermediate values.

NOTES - Include any comments not relevant in another section such as descriptions or unusual coding techniques, formulas, or release dependencies.

OPERATION - Include here a brief description, in general terms, of the logical steps performed by the module. The steps may be numbered so that branches and loops may be easily described.

The prologue is used for most of the system's modules. However, in some cases a module consists of a collection of relatively unrelated subroutines that have been grouped together but contain little common code. In this case, the prologue page contains only the MODULE NAME section and a list of the subroutines that the module contains. Preceding each subroutine, there is a prologue exactly as described above, except that the heading MODULE NAME is replaced by the heading SUBROUTINE NAME.

If CSECTS are large (more than 4096 bytes) it is preferable that code and referenced data reside in the

same page. The proper use of LTOrg statements will accomplish this.

EJECT  
 label CSECT  
 ENTRY Statements  
 EXTRN Statements  
 USING Statements  
 Source Code  
 Constants (See Note)  
 LTOrg (See Note)  
 Working Storage (See Note)  
 System DSECTS and EQUATES  
 END

MOD is the 3 character module name without the component code.

Label is the formal module name consisting of the component code and the Module Name.

Note: Where possible, these are in the same storage page with the instructions.



APPENDIX B: VM/370 CODING CONVENTIONS

1. **FORMAT:**

col. 1 - labels  
col. 10 - op code  
col. 16 - operands  
col. 31, 36, 41, etc. - comments (See Item 2.)

2. **COMMENT:**

Approximately 75 per cent of the source code contains comments. Sections of code performing distinct functions are separated from each other by a comment section.

3. **CONSTANTS:**

Constants follow the executable code and precede the copy files and/or macros which contain dsects or system equates. Constants are defined in a section followed by a section containing initialized working storage, followed by working storage. Each of these sections are identified by a comment. Where possible for a module that is greater than a page, constants and working storage are within the same page in which they are referenced.

4. No program modifies its own instructions during execution.

5. No program uses its own unlabeled instructions as data.

6. **REGISTER USAGE:** For CP, in general

<u>Register</u>	<u>Use</u>
6	- RCHBLOK, VCHBLCK
7	- RCUBLOK, VCUBLOK
8	- RDEVBLCK, VDEVBLCK
10	- IOBLOK
11	- VMBLOK
12	- Base register for modules called via SVC

- 13 - SAVEAREA for modules called via SVC
- 14 - Return linkage for modules called via BALR
- 15 - Base address for modules called via BALR

For Virtual to Real address translation:

- 1 - Virtual Address
- 2 - Real Address

7. When describing an area of storage in mainline code, a copy file, or a macro, DSECT is issued containing DS instructions.

8. Meaningful names are used instead of self-defining terms for example 5,X'02',C'I') to represent a quantity (for example absolute address, offset, length, register, etc.). All labels, displacements, and values are symbolic. All bits should be symbolic and defined by EQU. For example:

```
VMSTATUS EQU X'02'
```

To set a bit, use:

```
OI BYTE,BIT
```

Where BYTE = name of field, BIT is an EQU symbol.

To reset a bit, use:

```
NI BYTE,255-BIT
```

To set multiple bits, use:

```
IO BYTE,BIT1+BIT2
```

etc....

All registers are referred to as:

```
R0, R1, ....., R15.
```

All lengths of fields or blocks are symbolic, i.e. length of VMBLOK is:

```
VMBLOKSZ EQU *-VMBLOCK
```

9. Avoid absolute relative addressing in branches and data references, (that is, location counter value (\*) or symbolic label plus or minus a self-defining term used to form either a displacement or offset).
10. When using a single operation to reference multiple values, specify each value referenced, for example:

```
LM R2,R4,CONT SET R2=CON1
                SET R3=CON2
                SET R4=CON3
.
.
.
CON1 DC F'1'
CON2 DC F'2'
CON3 DC F'3'
```

11. Do not use PRINT NOGEN.
12. Module Names: Control Section Names and External References are as follows:  
  
Control Section or Module Name  
The first three letters of the name are the assigned component code.

Example: DMK

The next three letters of the Module Name identify the module and must be unique.

Example: DSP

This three letter unique module identifier is the label of the TITLE card.

Each entry point or external reference must be prefixed by the six letter unique identifier of the module.

Example: DMKDSPCH

13. TITLE Card:  
  
DSP TITLE 'DMKDSP VM/370 DISPATCHER VERSION 1 LEVEL 0'
14. PTF Card Example:  
  
CP/CMS: PUNCH 'xxxxxxxx APPLIED'  
  
Where xxxxxxxx = APAR Number Response
15. Error messages:  
  
There should not be any insertions into the message at execution time and the length of the message should be resolved by the assembler. If insertions must be made, the message must be assembled as different DC statements, and the insert positions are to be individually labeled.
16. For all RX instructions use ',' to specify the base register when indexing is not being used, i.e.

```
L R2,AB(,R4)
```

APPENDIX C. CP EQUATE SYMBOLS

CP DEVICE CLASSES, TYPES, MODELS AND FEATURES

CLASTERM	EQU	X'80'	Terminal Device Class
TYP2700	EQU	X'40'	2700 Bisync line
TYP2955	EQU	TYP2700	2955 Communications Line
TYPTLE2	EQU	X'20'	Telegraph Terminal Control Type II
TYPTTY	EQU	X'20'	Teletype Terminal
TYPIEM1	EQU	X'10'	IBM Terminal Control Type I
TYP2741	EQU	X'18'	2741 Communications Terminal
TYP1050	EQU	X'14'	1050 Communications Terminal
TYPUNDEF	EQU	X'1C'	Terminal device type is undefined
TYP3210	EQU	X'00'	3210 Console
TYP3215	EQU	TYP3210	3215 Console
TYP2150	EQU	TYP3210	2150 Console
TYP1052	EQU	TYP3210	1052 Console
CLASGRAF	EQU	X'40'	Graphics Device Class
TYP2250	EQU	X'80'	2250 Display Unit
TYP2260	EQU	X'40'	2260 Display Station
TYP2265	EQU	X'20'	2265 Display Station
TYP3066	EQU	X'10'	3066 Console
TYP1053	EQU	X'08'	1053 Printer
TYP3277	EQU	X'04'	3277 Display Station
TYP3284	EQU	X'02'	3284 Printer
TYP3286	EQU	TYP3284	3286 Printer
CLASURI	EQU	X'20'	Unit Record Input Device Class
TYPRDR	EQU	X'80'	Card Reader
TYP2501	EQU	X'81'	2501 Card Reader
TYP2540R	EQU	X'82'	2540 Card Reader
TYP3505	EQU	X'84'	3505 Card Reader
TYP1442R	EQU	X'88'	1442 Card Reader/Punch
TYP2520R	EQU	X'90'	2520 Card Reader/Punch
TYPTIMER	EQU	X'40'	Timer
TYPTR	EQU	X'20'	Tape Reader
TYP2495	EQU	X'21'	2495 Magnetic Tape Cartridge Reader
TYP2671	EQU	X'22'	2671 Paper Tape Reader
TYP1017	EQU	X'24'	1017 Paper Tape Reader
CLASURO	EQU	X'10'	Unit Record Output Device Class
TYPPUN	EQU	X'80'	Card Punch Device
TYP2540P	EQU	X'82'	2540 Card Punch
TYP3525	EQU	X'84'	3525 Card Punch
TYP1442P	EQU	X'88'	1442 Card Punch
TYP2520P	EQU	X'90'	2520 Card Punch

TYPVRT	EQU	X'40'	Printer
TYP1403	EQU	X'41'	1403 Printer
TYP3211	EQU	X'42'	3211 Printer
TYP1443	EQU	X'44'	1443 Printer
TYPTP	EQU	X'20'	Tape Punch
TYP1018	EQU	X'24'	1018 Paper Tape Punch
FTRUCS	EQU	X'01'	UCS Feature
CLASTAPE	EQU	X'08'	Magnetic Tape Device Class Units
TYP2401	EQU	X'80'	2401 Tape Drive
TYP2415	EQU	X'40'	2415 Tape Drive
TYP2420	EQU	X'20'	2420 Tape Drive
TYP3410	EQU	X'08'	3410 Tape Drive
TYP3420	EQU	X'10'	3420 Tape Drive
FTR7TRK	EQU	X'80'	7-Track Feature
FTRDLDNS	EQU	X'40'	Dual Density Feature
FTRTRANS	EQU	X'20'	Translate Feature
FTRDCONV	EQU	X'10'	Data Conversion Feature
CLASDASD	EQU	X'04'	Direct Access Storage Device Class
TYP2311	EQU	X'80'	2311 Disk Storage Drive
TYP2314	EQU	X'40'	2314 Disk Storage Facility
TYP2319	EQU	TYP2314	2319 Disk Storage Facility
TYP2321	EQU	X'20'	2321 Data Cell Drive
TYP3330	EQU	X'10'	3330 Disk Storage Facility
TYP2301	EQU	X'08'	2301 Parallel Drum
TYP2303	EQU	X'04'	2303 Serial Drum
TYP2305	EQU	X'02'	2305 Fixed Head Storage Device
FTR2311T	EQU	X'20'	(= VDEV231T) Top half of 2314 used as 2311
FTR2311B	EQU	X'10'	(= VDEV231B) Bottom Half of 2314 used as 2311
FTRRSRL	EQU	X'02'	Reserve/Release are valid CCW op codes
			(control unit has a 2-channel switch)
CLASSPEC	EQU	X'02'	Special Devices
TYPCTCA	EQU	X'80'	Channel to Channel Adapter

## MACHINE USAGE

### Bits defined in standard extended PSW

EXTMCODE EQU	X'08'	Bit 12 - Extended Mode
MCHEK EQU	X'04'	Bit 13 - Machine check enabled
WAIT EQU	X'02'	Bit 14 - Wait state
PROEMODE EQU	X'01'	Bit 15 - Problem state

### Bits defined in extended PSW

PERMODE EQU	X'40'	Bit 01 - PER enabled
MODE31 EQU	X'08'	Bit 04 - 31 bit mode addressing
TRANMODE EQU	X'04'	Bit 05 - Translate mode
IOMASK EQU	X'02'	Bit 06 - Summary I/C Mask
EXTMASK EQU	X'01'	Bit 07 - Summary external mask

### Bits defined in channel status word - CSW

ATTN EQU	X'80'	Bit 32 - Attention
SM EQU	X'40'	Bit 33 - Status modifier
CUE EQU	X'20'	Bit 34 - Control unit end
BUSY EQU	X'10'	Bit 35 - Busy
CE EQU	X'08'	Bit 36 - Channel end
DE EQU	X'04'	Bit 37 - Device end
UC EQU	X'02'	Bit 38 - Unit check
UE EQU	X'01'	Bit 39 - Unit exception
PCI EQU	X'80'	Bit 40 - Program-control interrupt
IL EQU	X'40'	Bit 41 - Incorrect length
PRGC EQU	X'20'	Bit 42 - Program check
PRTC EQU	X'10'	Bit 43 - Protection check
CDC EQU	X'08'	Bit 44 - Channel data check
CCC EQU	X'04'	Bit 45 - Channel control check
IFCC EQU	X'02'	Bit 46 - Interface control check
CHC EQU	X'01'	Bit 47 - Chaining check

### Bits defined in channel command word - CCW

CD EQU	X'80'	Bit 32 - Chain data
CC EQU	X'40'	Bit 33 - Command chain
SILI EQU	X'20'	Bit 34 - Suppress incorrect length indication
SKIP EQU	X'10'	Bit 35 - Suppress data transfer
PCIF EQU	X'08'	Bit 36 - Program-control interrupt fetch
IDA EQU	X'04'	Bit 37 - Indirect data address

### Bits defined in sense byte 0 -- Common to most devices

CMDREJ EQU	X'80'	Bit 0 - Command reject
INTREQ EQU	X'40'	Bit 1 - Intervention required
BUSCUI EQU	X'20'	Bit 2 - BUS out
EQCHK EQU	X'10'	Bit 3 - Equipment check
DATACHK EQU	X'08'	Bit 4 - Data check

EXTENDED CONTROL REGISTERS

## Bits defined in Control register 0

		Byte 0	
BLKMPX	EQU	X'80'	Bit 00 - Enable block multiplexing
SSMSUPP	EQU	X'40'	Bit 01 - Enable SSM suppression
		Byte 1	
PAGE4K	EQU	X'80'	Bit 08 - Use 4K pages
PAGE2K	EQU	X'40'	Bit 09 - Use 2K pages
SEG1M	EQU	X'08'	Bit 12 - Use 1M segments
		Byte 2	
CKCMASK	EQU	X'08'	Bit 20 - Mask on clock comparator intercept
CPTMASK	EQU	X'04'	Bit 21 - Mask on CPU timer intercept
		Byte 3	
INTMASK	EQU	X'80'	Bit 24 - Mask on interval timer intercept
KEYMASK	EQU	X'40'	Bit 25 - Mask on operator key intercept
SIGMASK	EQU	X'20'	Bit 26 - Mask on external signals 2-7

## Bits defined in Control register 9

		Byte 0	
PERSUBR	EQU	X'80'	Bit 00 - Monitor successful branches
PERIFET	EQU	X'40'	Bit 01 - Monitor instruction fetches
PERSALT	EQU	X'20'	Bit 02 - Monitor storage alteration
PERGPRS	EQU	X'10'	Bit 03 - Monitor register alteration

## Bits defined in Control register 14

		Byte 0	
HARDSIOP	EQU	X'80'	Bit 00 - Check stop control
SYNLOG	EQU	X'40'	Bit 01 - Synchronous logout control
IOLOG	EQU	X'20'	Bit 02 - I/O logout control
RECOVERPT	EQU	X'08'	Bit 04 - Recovery report mask
CONFGRPT	EQU	X'04'	Bit 05 - Configuration report mask
DAMAGRPT	EQU	X'02'	Bit 06 - External damage report mask
WARNGRPT	EQU	X'01'	Bit 07 - Warning condition report mask
		Byte 1	
ASYNELOG	EQU	X'80'	Bit 08 - Asynchronous extended logout control
ASYNFLOG	EQU	X'40'	Bit 09 - Asynchronous fixed logout control

## CONTROL PROGRAM USAGE

### Bits defined for TRANS macro

BRING	EQU	X'80'	Bring requested page
DEFER	EQU	X'40'	Defer execution until page in storage
LOCK	EQU	X'20'	Lock page for I/O operation
IOERFIN	EQU	X'10'	Return I/O errors to caller
SYSTEM	EQU	X'08'	Call to DMKPTRAN for system virtual machine space

### Bits defined for terminal I/O

ERRMSG	EQU	X'04'*256	Control program error message
PRIORITY	EQU	X'02'*256	Queue and start this message immediately
VMGENIO	EQU	X'01'*256	Virtual machine generated I/O request
LOGDROP	EQU	X'80'	LOGOUT and drop line after output message
LOGHCLD	EQU	X'40'	LOGCUT & hold line after output message
NORET	EQU	X'20'	Return immediately after call
DFRET	EQU	X'10'	FRET Buffer after write
NOAUTO	EQU	X'08'	No automatic carriage return
EDIT	EQU	X'08'	Edit input for corrections
ALARM	EQU	X'04'	Sound the alarm
UCASE	EQU	X'04'	Translate input to upper case
OPERATOR	EQU	X'02'	Message for operator
NOTIME	EQU	X'01'	Do not time stamp message





APPENDIX D. DASD RECORD FORMATS

RECORD 0, 8 BYTES (PAGE BIT MAP)

Used to flag pages that are in use or have bad recording area. Devices that do not use all 64 bits (64 pages per cylinder) have the unused bits turned on.

Examples

RECORD 0 CYLINDER 0 Only

32 Pages/cylinder 2314,2319

```
*|E0 00 00 00 FF FF FF FF|
```

|  
11100000

57 pages/cylinder 3330

```
|E0 00 00 00 00 00 00 7F|
```

24 pages/cylinder 2305

```
|E0 00 00 FF FF FF FF FF|
```

\* The first three pages of cylinder 0 are always flagged in use, since they are used by CP. On all other cylinders, the first byte hex '00' unless the disk area is flagged bad. Record 0 of all tracks other than track 0 is initialized to hex '00'.

All Page Records, 4096 Bytes Each

2314 and 2319 32 pages/cylinder  
3330 series 57 pages/cylinder  
2305 24 pages/cylinder

Cylinder 0 contains less pages because this area is used by CP.

RECORD 1 (24 BYTES)

IPL record - Puts system into wait state if storage device is IPLed.

```
|00020000 0000000C 03000000 20000000 00000000 00000000|
```

RECORD 2, 4096 BYTES

Check point record - this is the CHECKPOINT program load at CP IPL time to retrieve and save control information for a warm start.

RECORD 3

4 byte key of VOL1  
80 byte data record

Key

VOL1

Record

Bytes	
1-20	E5D6D3F1 xx----->xxF000 00000005 0000C000
21-40	0040----->40
41-60	4000----->00C3D7 F3F7F040 40404040 40--->40
61-80	40----->40

Where:

xx->xx is a 6 byte label

Bytes 13-16 is a pointer to the VTOC

Bytes 46-50 identify the system

Bytes 52-55 is a pointer to the active directory

RECORD 4

1024 bytes Track 0 Cylinder 0

Allocation byte map - used to identify Cylinder 1 usage. Each byte identifies one cylinder.

C -> all 0<-

00000100 040200----->FF 0000----->0000
--

\*

\* FF defines the last cylinder + 1 that can be allocated. This varies depending on the device.

00 = temporary  
01 = permanent  
02 = T-disk  
04 = directory

RECORD 5

44 bytes key Track 0, Cylinder 0 96 bytes data area

Format 4 OS DSCB type label - used to be compatible with OS.

04----->04

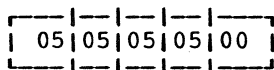
44 Key

FORMAT 4 LABEL

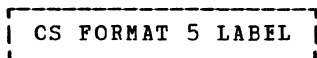
96 Byte Data

RECORD 6

44 bytes key Track 0, Cylinder 0 96 bytes data area  
 Format 5 OS DSCB type lable for ccmpatibility with OS.



44 Byte Key



96 Byte Data Area

RECORD F3

4096 bytes - 1 page, track 0 or track 1  
 F3 Record is reserved for CPsystem use. Referred to as filler record.

RECORD F4

1624 bytes, Track 1 (2314, 2319 only)  
 F4 used only on 2314 and 2319 devices to align Record 4 in proper position on track.

RECORD 4

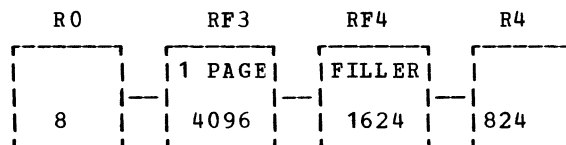
824 bytes track 1, cylinder 0 =2314, 2319 only)  
 First segment of Record 4 to be used for paging.

2314 RECORD LAYOUT

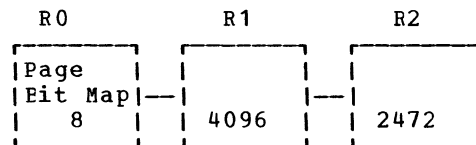
CYLINDER 0, TRACK 0

R0	R1	R2	Key	R3	R4	Key	R5	Key	R6
Page	I	Check	V	VOL1	Alloc	Format	Format	Format	
Bit	P	Point	O	Label	Byte	4	5		
Map	L		L		Map				
8	24	4096	4	80	1024	44	96	44	96

Cylinder 0, Track 1

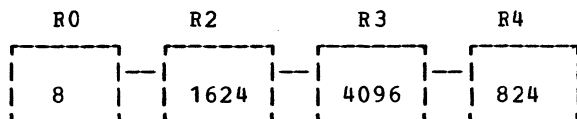


ALL CYLINDERS EXCEPT 0, TRACK 0

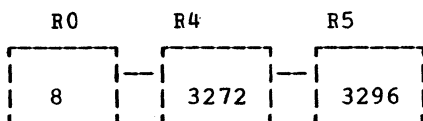


These records appear as above formats if cylinder is 0.

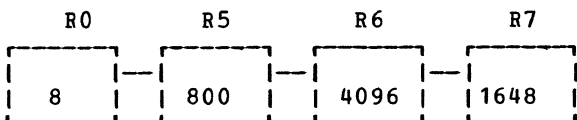
Track 1



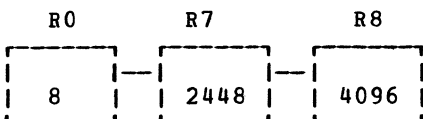
Track 2



Track 3



Track 4



Note: Track 0 to 4 are repeated for tracks 5 to 9 (=R9\_R16), 10 to 14, (=R17\_R24), and 15 to 19 (=R25\_R32). The last record is R32.

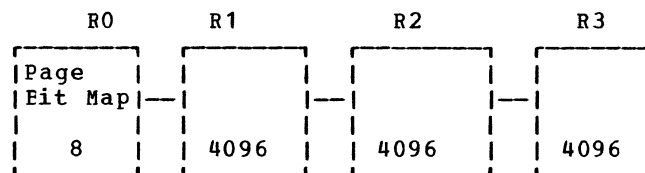
3330 SERIES RECORD LAYOUT

CYLINDER 0, TRACK 0

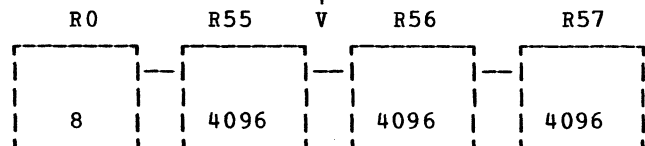
R0	R1	R2	Key	R3	R4	Key	R5	Key	R6	RF3
Page	I	Check	V	VOL1	Byte	Format	Format	1		
Bit	P	Point	O	Label	Map	4	5	Page		
Map	L		L							
8	24	4096	4	80	1024	44	96	44	96	4096

ANY CYLINDER EXCEPT 0

Track 0



Track 18



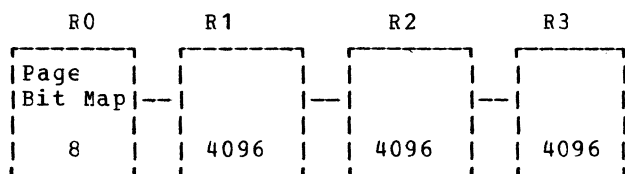
2305 MODEL 1 and MODEL 2

CYLINDER 0, TRACK 0

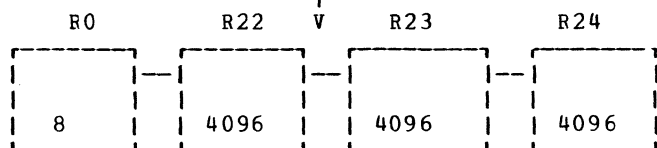
R0	R1	R2	Key	R3	R4	Key	R5	Key	R6	RF3
Page	I	Check	V	VCL1	Byte		Format		Format	1
Bit	P	Point	O	Label	Map		4		5	Page
Map	L		L							
8	24	4096	4	80	1024	44	96	44	96	4096

ANY CYLINDER EXCEPT 0

Track 0



Track 7





*Ken Jones*

This Newsletter No. SN20-2624  
Date August 15, 1973  
Base Publication No. SY20-0880-1  
File No. S370-37 (VM/370  
Release 1 PLC 9)  
Previous Newsletters None

## IBM Virtual Machine Facility/370: Control Program (CP) Program Logic

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This Technical Newsletter, a part of Release 1 PLC 9 of IBM Virtual Machine Facility/370 provides replacement pages for your publication. These replacement pages remain in effect for subsequent VM/370 releases unless specifically altered. Pages to be removed and/or inserted are listed below.

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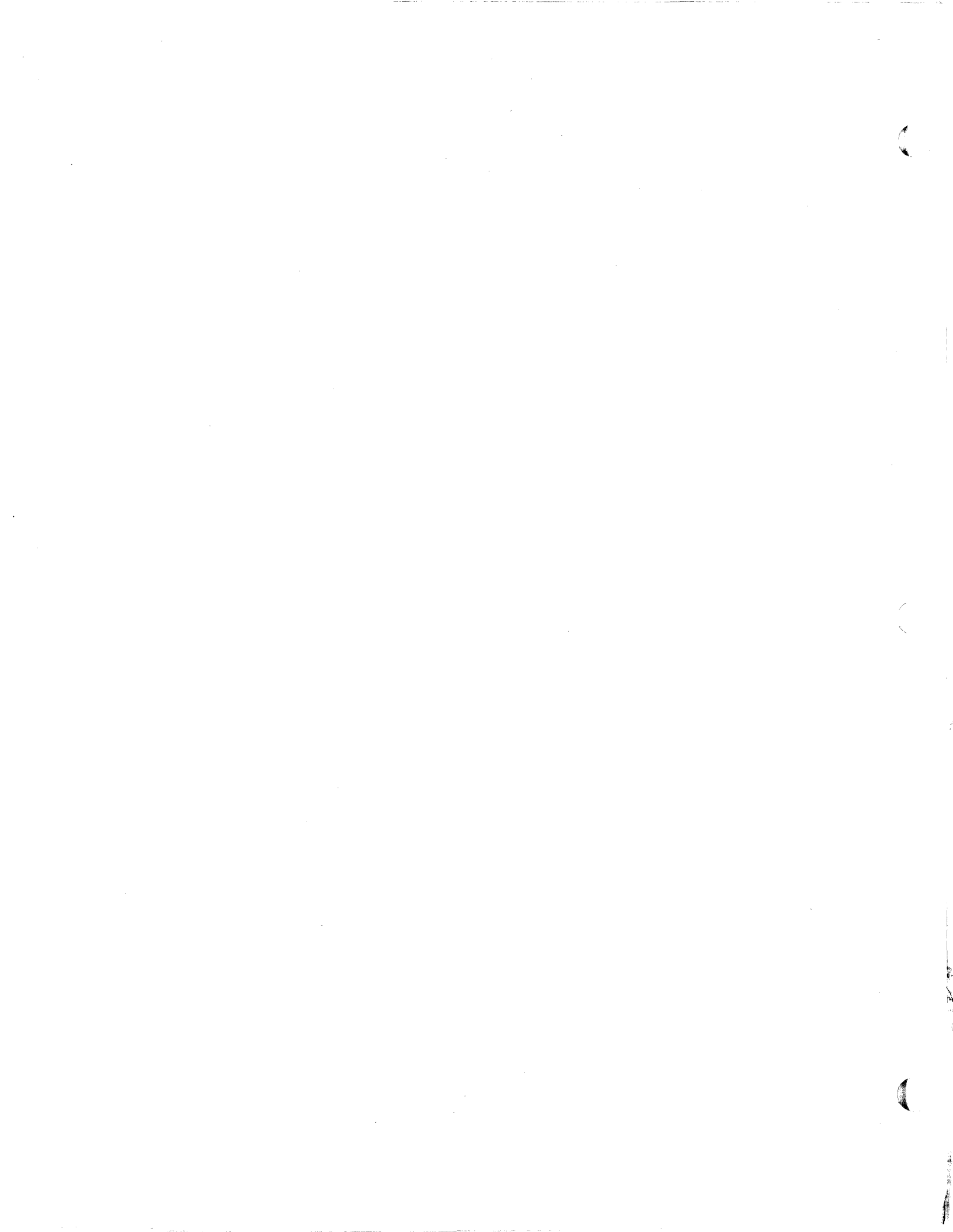
### Summary of Amendments

This Technical Newsletter contains changes that reflect VM/370 support for:

- The IBM System/370 Model 168
- The IBM 2860 Selector Channel
- The IBM 2870 Multiplexor Channel
- The IBM 2880 Block Multiplexor Channel
- The IBM 2305 Fixed Head Storage, Model 1
- User Accounting Option
- Virtual Console Spooling

See the Summary of Amendments page for further details.

**Note:** Please file this cover letter at the back of your publication to provide a record of changes.





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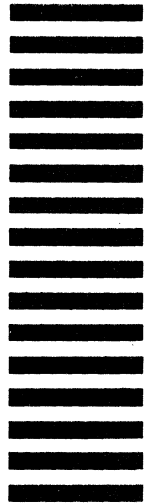
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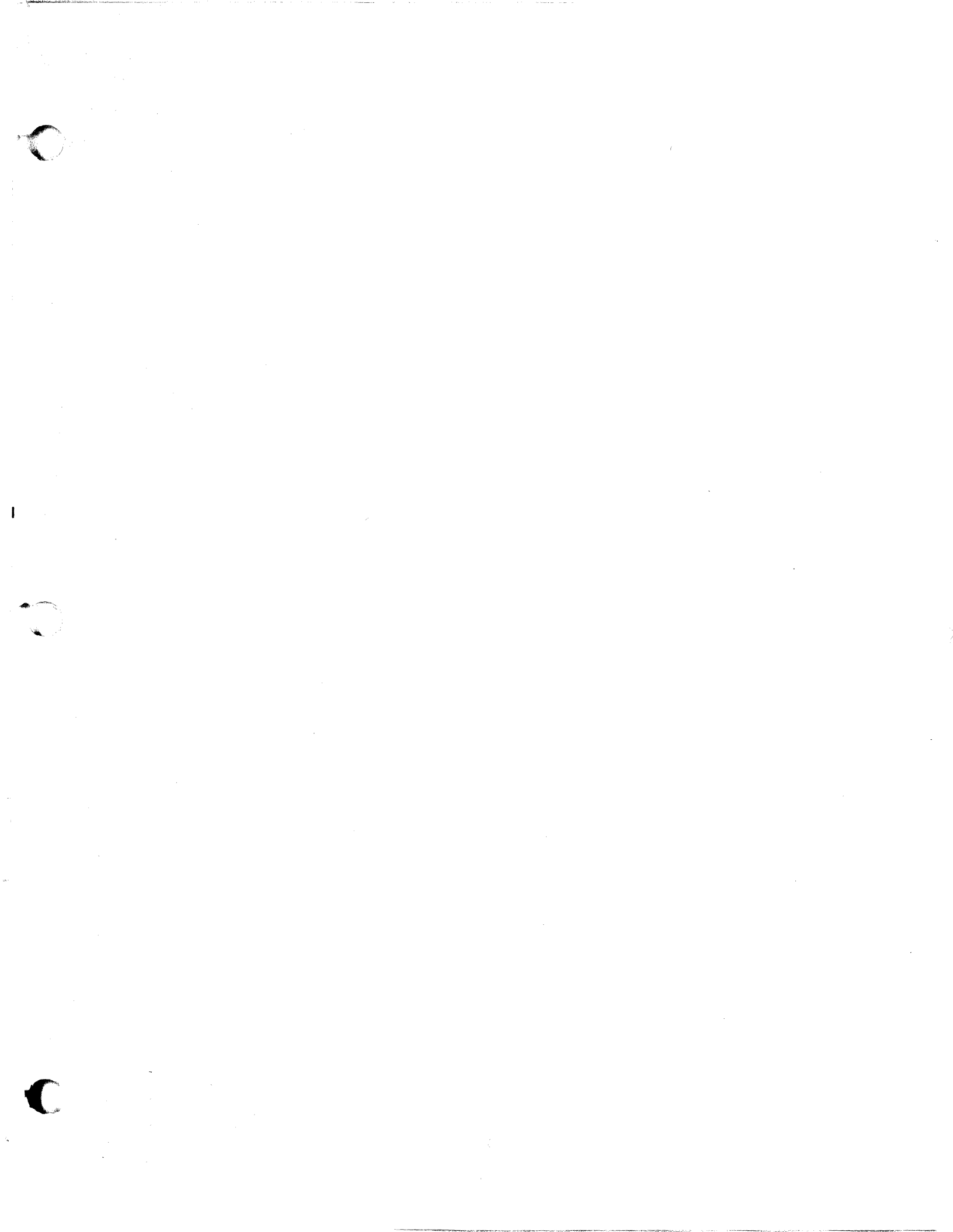
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