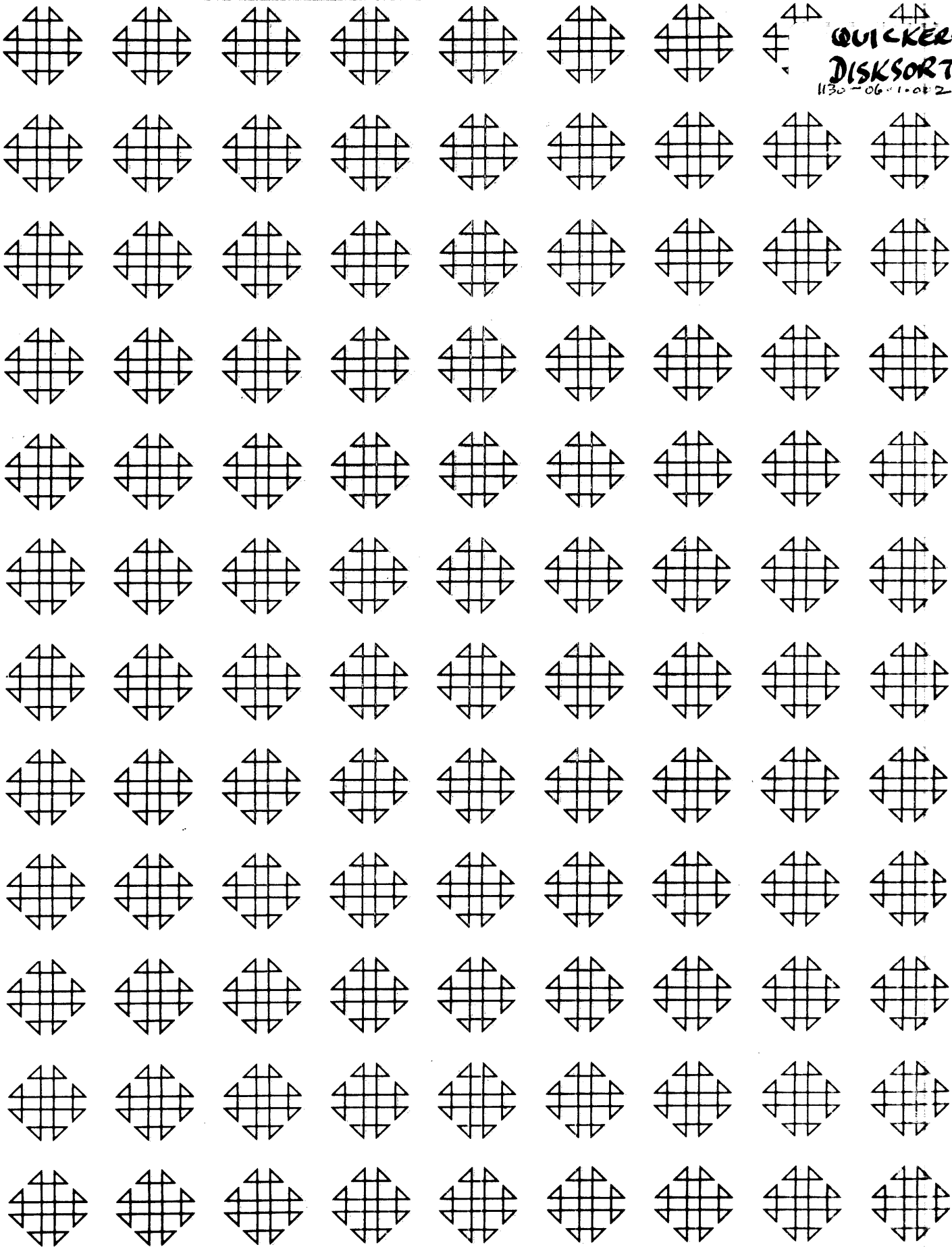
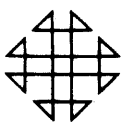


Sort Subprograms for IBM 1130, Quickersort,  
Disksort 1 and Disksort 2

1130-06.1.002

**QUICKER  
DISKSORT**  
1130-06.1.002



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Sort Subprograms for the IBM 1130;  
Quickersort, Disksort 1 and Disksort 2

Robert N. Kubik  
The Babcock-Wilcox Company  
1201 Kemper Street  
Lynchburg, Virginia

March 21, 1967

Modifications or revisions to this program, as they occur, will be announced in the appropriate Catalog of Programs for IBM Data Processing Systems. When such an announcement occurs, users should order a complete new program from the Program Information Department.

PROGRAM ABSTRACT

1. TITLE (If subroutine, state in Title): QSORT, DSRT 1, DSRT 2 ( 3 subroutines)  
Subject Classification: 06.1
2. Author; Organization: Mr. Robert N. Kubik, The Babcock-Wilcox Company  
Date: 3/26/67 Users Group Membership Code: \_\_\_\_\_
3. Direct Inquiries to Name: Mr. Robert N. Kubik, The Babcock-Wilcox Company  
1201 Kemper Street, Lynchburg, Virginia Phone: 703 846-7371
4. Description/Purpose: (5. Method: 6. Restriction/Range; When Applicable):  
Three sort subprograms written in FORTRAN for the IBM 1130 are included. First, sort based on Algorithm 271 as published in the ACM Communications requires only a minimum 1130. Second, a disk sort which requires 8K memory and disk. Third, a disk sort which requires 4K memory and disk but is less effecient. These would not be practical for sorting large files. These subprograms are available at this time for the 1130.
7. Specifications (Check or fill in appropriate spaces):
- a. Storage used by Program: 4K, 8K, 4K, respectively
- b. Equipment required by program: Card \_\_\_\_\_; Magnetic Tape \_\_\_\_\_; Number of Drives \_\_\_\_\_; Paper Tape \_\_\_\_\_; Disk File 2&3; Number of Drives \_\_\_\_\_; TNS, TNF, MF, \_\_\_\_\_; Auto divide \_\_\_\_\_; Indirect addressing \_\_\_\_\_; Floating Point Hardware \_\_\_\_\_; 1620 Model I \_\_\_\_\_; Model II \_\_\_\_\_; 1443 Printer \_\_\_\_\_; Index Registers \_\_\_\_\_; Binary Capabilities \_\_\_\_\_; other (specify) \_\_\_\_\_
- Can program be used on lesser machine? NO. Specify which requirements can be easily removed \_\_\_\_\_
- c. Programmed in: Fortran without Format \_\_\_\_\_; Fortran with Format \_\_\_\_\_; Fortran II \_\_\_\_\_; Other Fortran (specify) \_\_\_\_\_; SPS (specify assembler used) \_\_\_\_\_; Other (specify) 1130 Fortran
- d. Type of Program: Mainline, Complete \_\_\_\_\_; Subroutine X; If subroutine, for use with SPS (specify type of SPS) \_\_\_\_\_; Fortran (specify type of Fortran) \_\_\_\_\_; Other (specify) \_\_\_\_\_
8. Additional Remarks: \_\_\_\_\_  
\_\_\_\_\_  
\_\_\_\_\_  
\_\_\_\_\_  
\_\_\_\_\_

1A

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DECK KEY

Deck #

1	FORTRAN	deck	QSORT	cards 1 - 110
2	FORTRAN	deck	DSRT1	cards 1 - 168
3	FORTRAN	deck	MAINL	cards 1 - 30
4	Sample problem input			cards 1 - 7
5	FORTRAN	deck	DSRT2	cards 1 - 94

card sequence numbers are in columns 76 - 80

Program Abstract

Three sort subprograms are included in this report. A FORTRAN version of Algorithm 271 Quicksort, published in CACM, vol. 8, No. 11, Oct., 1965, by R. S. Scowen; which is a very fast internal sort for randomly distributed keys. It is herein called QSORT. It requires minimal storage, i.e. storage for the records to be sorted plus one record. Run time in favorable cases can be approximated by:

$$\left( \frac{\text{number of records}}{80} \right)^{1.3} = \text{run time in seconds,}$$

for instance 1200 records of two words each were sorted in 30.4 seconds. However, unfavorable key distribution can result in substantially longer run time. Cases have been observed where the exponent in the above formula approaches 1.9.

DSRT1 is a disk sort for files exceeding core storage available. It uses QSORT to create ordered strings of data. Then a series of passes are made which can best be described as combining the logic of pair exchange and merging. DSRT1 requires 2560 plus the number of words in a record, words of working storage. Run time is approximately

$$\frac{(\text{number of records})^2}{10,000} \times \text{number of words/record} = \text{run time in seconds.}$$

DSRT2 is a disk sort for files exceeding fast storage size available. It is a direct application of the same algorithm as QSORT. It requires 1632 words working storage, thereby minimizing storage requirement. In all cases tested to date it has proven substantially slower than DSRT1. It is possible, however, it would prove to be faster than DSRT1 on relatively large files with large records, (80 words per record or more).

These sorts are intended for sorting keys and tags (e.g. record numbers) so that one can find records in a larger file in proper sequence for updating or report preparation. They can, however, be used as any sort routines would be - provided the records have less than 321 words/record; the keys are integers and are grouped in major to minor sequence as the first items of each record. Suggestions are included for circumventing this restriction on the keys and for speeding up these routines. These subroutines should be suitable for the IBM 1800 although they have not been tested on that machine.

March 21, 1967

This program and its documentation were written by Mr. Robert N. Kubik, ( the Babcock-Wilcox Company, 1201 Kemper Street, Lynchburg, Virginia, 703 846-7371). They have been submitted to the Program Information Department for general distribution in the expectation that they may prove useful to other members of the data processing community. The program and its documentation are, essentially, in the author's original form and have not been subject to any formal testing. IBM only serves as the distribution agency in supplying this program. It is the user's responsibility to determine the usefulness and technical accuracy of the program in his own environment. This program is not part of the Programming System (Type I) and Application Programs (Type II). Questions concerning the use of the program should be directed to the author. Any changes to the program will be reflected in the appropriate catalog of programs; however, the changes will not be distributed automatically to users.

-5-

## USER INFORMATION

Purpose: QSORT is an internal sort for files which can be held in core storage. DSRT1 and DSRT2 are for files exceeding the size of available core storage.

Advantages: QSORT minimizes working storage and is very fast for randomly distributed keys. DSRT1 is relatively efficient. DSRT2 minimizes working storage requirement. Sorting large files is awkward at best on a one disk system. Usually the most efficient solution is to sort only keys and tags (e.g. record numbers) and then access the file in its existing order by means of the table so generated. These subroutines appear to be adequate for that purpose.

Method: QSORT and DSRT2 continually split the file into parts such that the keys of all records in one part are less than the keys of all records in the other part, with a third part in the middle consisting of a single record. A record with key  $t$  is chosen arbitrarily,  $i$  and  $j$  are the lower and upper records of the segment being split. After the split has taken place the  $q$ th record will have been found such that the key of the  $q$ th record is  $t$  and the key of the  $i$ th record  $\leq t \leq$  the key of the  $j$ th record, for all  $i$  and  $j$  such that  $i \leq I < q < J \leq j$ . The program then performs operations on the two segments; the  $i$ th thru  $(q-1)$ th records and the  $(q+1)$ th thru  $j$ th records as follows. The smaller segment is split and the position of the larger segment is stored in the  $l$  (lower temporary) and  $u$  (upper temporary) arrays. If the segment to be split has two or fewer elements it is sorted and another segment obtained from the  $l$  and  $u$  arrays. When no more segments remain, the array is completely sorted.

DSRT1 first sorts groups of records of approximate cylinder size by use of QSORT. It then handles the file in groups of records of approximately one half cylinder size. It makes successive merge passes merging the second and third group, the fourth and fifth group and so on; and then merges the first and second group, the third and fourth and so on. These merge passes are alternated until the file is in order.

### Restrictions and range:

- number of records  $\leq 32,767$
- number of words/record  $\leq 320$
- keys numeric, the first words in the record and in major to minor sequence
- sorts file 1 as defined in FORTRAN DEFINE FILE statement
- integers are considered to be one word each

- 6 -

-in QSORT and DSRT2 p is the position of an arbitrary record in the file. The best value of p is one which splits the segment into two halves of equal size, thus if the file (segment) is roughly sorted, the middle element is an excellent choice. If the file (segment) is completely random the middle element is as good as any other. If, however, the file (segment) is such that the parts i...p and p+1..j are both sorted the middle element could be very bad. Accordingly in some circumstances  $p = (i + j)/2$  should be replaced by  $p = (i + 3 \times j)/4$  or by a p selected at random. The method of selecting p used in these routines will prove to be best in the majority of cases, however.

Program Requirements: These subroutines as distributed are designed for use within FORTRAN programs operated under the 1130 Disk Monitor System. QSORT requires a 4K machine only, although the more space that is available the larger the file that can be sorted. DSRT1 requires 8K storage and a disk. DSRT2 requires 4K storage and a disk.

Timing: In QSORT the number of words in the key influence run time very little; the number of words in a record somewhat more, but not substantially. The largest effects on run time are the number of records and the distribution of keys. A number of cases run with random key distribution have exhibited a run time characterized by the expression:

$$\left( \frac{\text{number of records}}{80} \right)^{1.3} = \text{run time in seconds,}$$

typical of these cases are the following:

number of records	words/record	words/key	time
300	2	2	5.4 sec.
600	2	2	12.3 sec.
1200	2	2	30.4 sec.

Under certain circumstances run time can be substantially longer (see Restrictions and Range) in fact series of cases have been observed which fit the expression:

$$\left( \frac{\text{number of records}}{80} \right)^{1.9} = \text{run time in seconds.}$$

Run time for larger records (say eight words or more) will be underestimated by this formula.

DSRT1 timings can be approximated by:

$$\frac{(\text{number of records})^2 \cdot \text{number of words per record}}{10,000} = \text{run time in seconds.}$$

Run time for large records (say eighty words or more) will be underestimated by this formula.

DSRT2 took longer on all cases run and no formula for predicting time has been derived. It is conjectured however, that on large files with 80 or more words per record it may prove more efficient than DSRT1. Some cases which have been timed follow.

records	words/record	DSRT1	DSRT2
128	80	202 sec.	230 sec.
250	20	120 sec.	315 sec.
1000	7	730 sec.	



### Program Modification Aids

All coding which compares keys has been surrounded by comment cards with \*\*\*\*\* on them and text explaining "insert your own code for non numeric keys or keys which are not the initial words of the record - from here .....- to here."

An assembly language version of QSORT would be worthwhile and should improve the speed of that routine by a factor of two. The algorithm should be clear from the FORTRAN code listing and flow chart.

DSRT1 could be speeded up for sorting small records by replacing the FORTRAN read/write statements by a routine which reads and writes blocks of records and delivers these to the DSRT1 subprogram on call. This would, however, be no advantage for large records and would impose additional restrictions on the program which creates the file to be sorted.

DSRT1 and DSRT2 sort FILE1. This could be changed by changing the file number in the disk read/write statement.

### Input-Output Description

QSORT is called by -

CALL QSORT (NREC, NWREC, NWKEY, A, T)

where NREC is the number of records

NWREC is the number of words/record

NWKEY is the number of words/key

A is a one dimensional array containing the file to be sorted

T is a working storage array of NWREC words.

It leaves the first NREC records in A sorted in ascending sequence.

DSRT1 is called by -

CALL DSRT1 (NRFIL, NWREC, NWKEY, ARRAY, TEMP)

where NRFIL is the number of records in the file

NWREC is the number of words/record

NWKEY is the number of words/key

ARRAY is a 2560 word array for working storage

TEMP is a working storage array of NWREC words

it leaves the first NRFIL records in FILE 1 sorted in ascending sequence.

DSRT2 is called by -

CALL DSRT2 (NREC, NWREC, NWKEY)

where NREC is the number of records in the file

NWREC is the number of words/record

NWKEY is the number of words/key

it leaves the first NREC records in FILE 1 sorted in ascending sequence.

### Sample Problem

A driver program called MAINL is included with the deck distributed. It reads test data, creates a file and prints it, calls DSRT1 (which calls QSORT) and prints the sorted file. This same driver may be used for DSRT2 by replacing the integer statement with INTEGER ARRAY (2560) and the call statement with CALL DSRT2 (NRFIL, NWREC, NWKEY)

### Operating Instructions

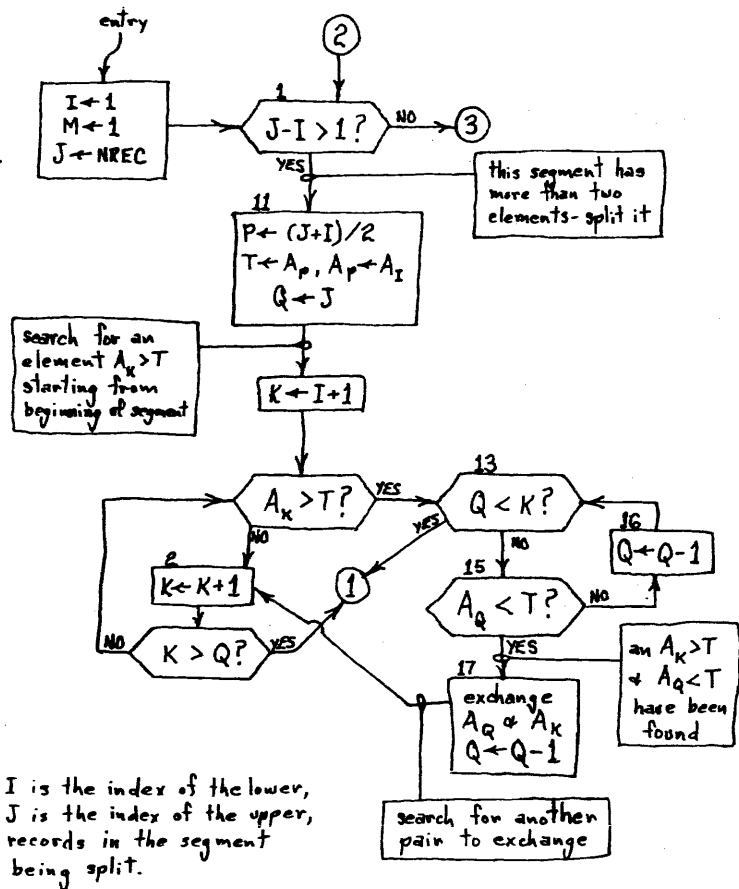
Illogical branches in DSRT1 result in a stop at PAUSE 1. If this occurs something is wrong with the machine or the program.

### References

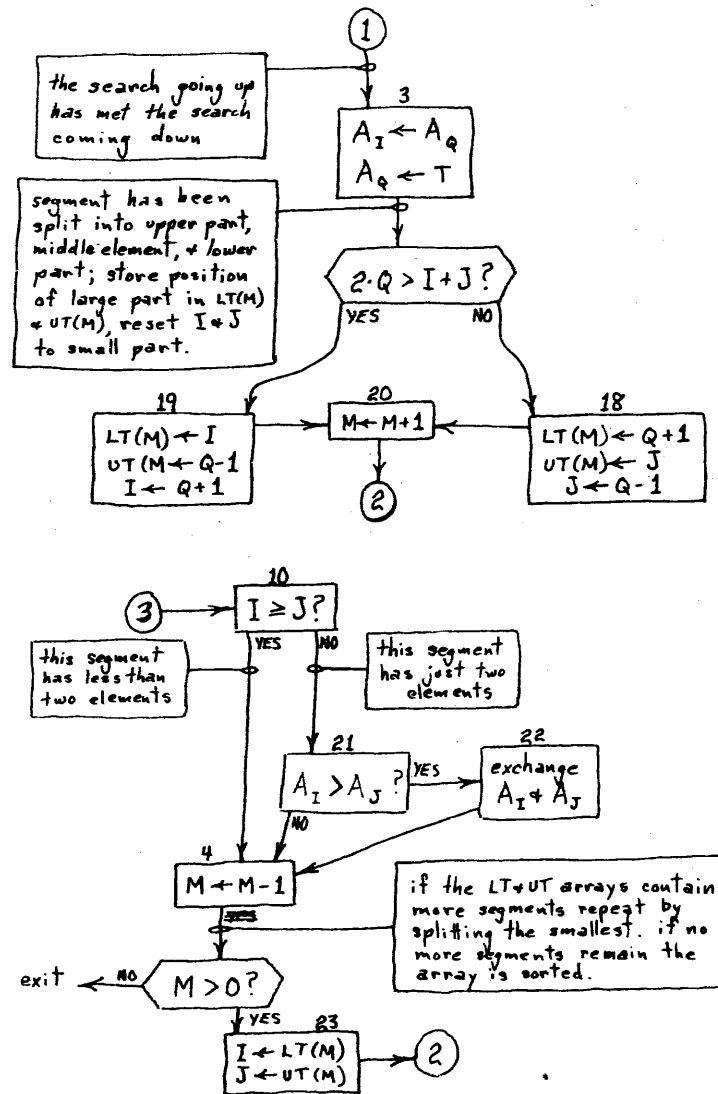
- C20-1639-0 Sorting Techniques, IBM Data Processing Techniques.
- Hoare, C. A. R. Algorithms 63 and 64. Comm. ACM 4 (July 1961), 321.
- Hibbard, Thomas N. Some combinatorial properties of certain trees with applications to searching and sorting. J. ACM 9 (Jan. 1962), 13.
- Hibbard, Thomas N. An empirical study of minimal storage sorting. Comm. ACM 6 (May 1963), 206-213.
- Scowen, R. S. Algorithm 271 (QUICKERSORT) Comm. ACM 11 (Nov. 1965), 669-670.

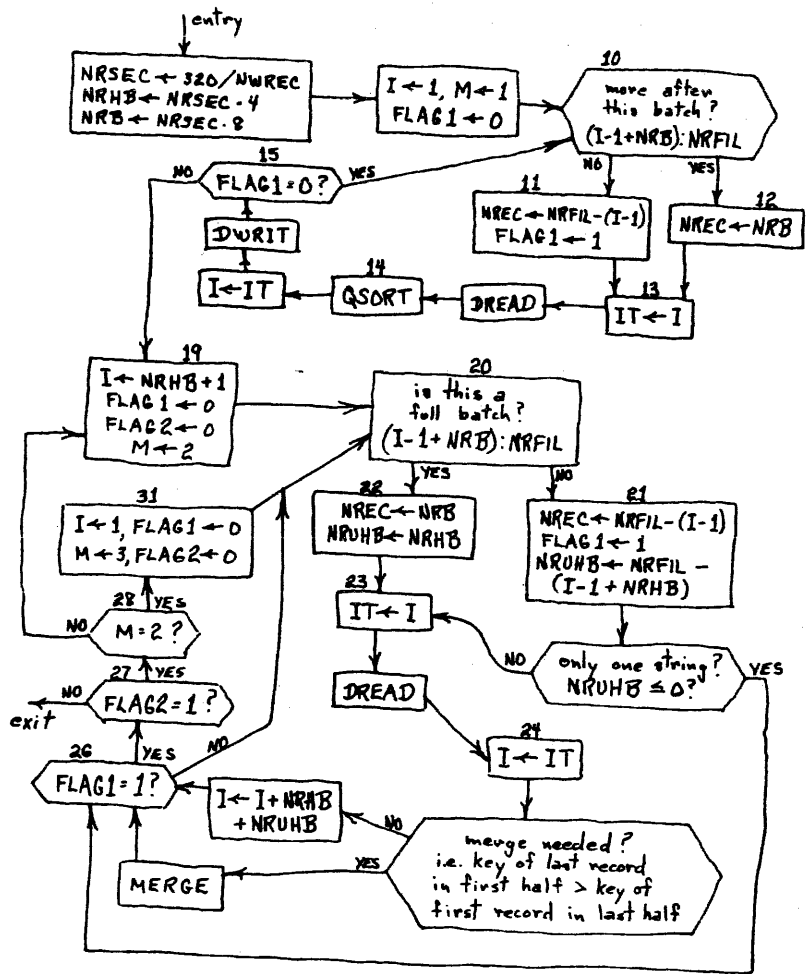
Systems Material

Statement numbers from the FORTRAN code appear above the appropriate box in these flowcharts as a cross index with the code. QSORT and DSRT2 have the same logic and the interpretation of that flowchart for either should pose no difficulty.

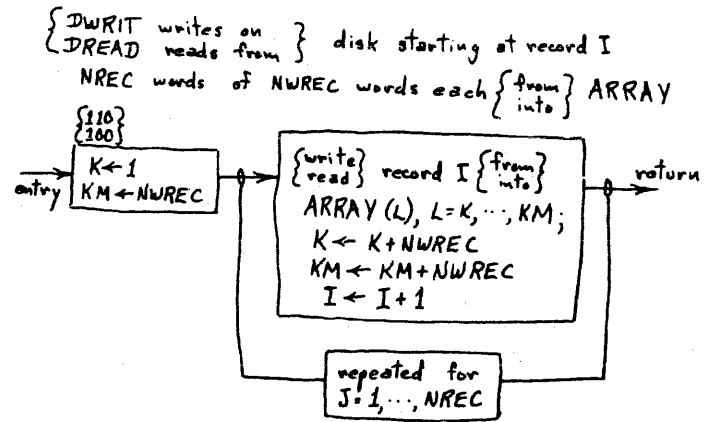


QSORT and DSRT2 Flowchart

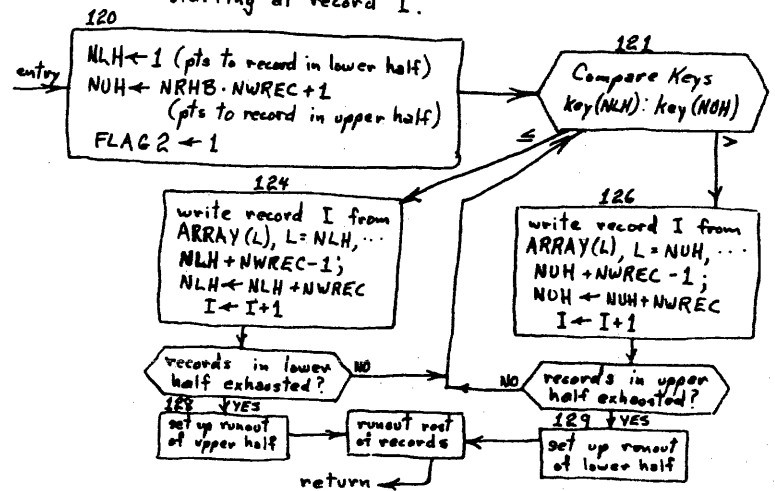




DSRT1 Flowchart



**DWRIT** writes on } disk starting at record I  
**DREAD** reads from }  
 NREC words of NWREC words each { from } ARRAY  
 into }



# Q SORT, DSRTJ, MAINL, complete edit run

// JOB T  
// FOR

0J01

\* NAME QSORT  
\* ONE WORD INTEGERS  
\* LIST ALL

PAGE 01

0J02  
0J03  
0J04

	SUBROUTINE QSORT (NREC,NWREC,NWKEY,A,T)	
C	QUICKSORT ALGORITHM 271 CACM VOL B NO 11	0J05
C	NREC IS THE NUMBER OF RECORDS, NWREC IS THE NUMBER	0J06
C	OF WORDS PER RECORD, NWKEY IS THE NUMBER OF WORDS	0J07
C	PER KEY, A IS THE ARRAY OF RECORDS WHICH TAKES	0J08
C	NREC*NWREC WORDS, T IS WORKING STORAGE OF NWREC WORDS	0J09
C	NREC*NWREC WORDS, T IS WORKING STORAGE OF NWREC WORDS	0J10
C	INTEGER A(32767), UT(15),LT(15),T(16383),Q,P,X	0J11
	M=1	0J12
	I=1	0J13
	J=NREC	0J14
	N	0J15
C	1 IF (J-I)10,10,11	0J16
	11 P=(J+I)/2	0J17
	NP=(P-1)*NWREC+1	0J18
	NI=(I-1)*NWREC+1	0J19
	DO 110 N=1,NWREC	0J20
	T(N)=A(NP)	0J21
	A(NP)=A(NI)	0J22
	NP=NP+1	0J23
	110 NI=NI+1	0J24
	Q=J	0J25
	I1=I+1	0J26
	DO 2 K=I1,Q	0J27
	C****INSERT YOUR OWN CODE FOR NON NUMERIC KEYS	0J28
	C****OR KEYS WHICH ARE NOT THE INITIAL WORDS	0J29
	C****OF THE RECORD -FROM HERE	0J30
	NK=(K-1)*NWREC+1	0J31
	DO 130 N=1,NWREC	0J32
	IF (A(NK)-T(N))2,130,13	0J33
	130 NK=NK+1	0J34
	GO TO 2	0J35
	C**** -TO HERE	0J36
	13 IF (Q-K)14,15,15	0J37
	C****INSERT YOUR OWN CODE FOR NON NUMERIC KEYS	0J38
	C****OR KEYS WHICH ARE NOT THE INITIAL WORDS	0J39
	C****OF THE RECORD -FROM HERE	0J40
	15 NQ=(Q-1)*NWREC+1	0J41
	DO 150 N=1,NWKEY	0J42
	IF (A(NQ)-T(N)) 17,150,16	0J43
	150 NQ=NQ+1	0J44
	GO TO 16	0J45
	C**** -TO HERE	0J46
	17 NK=(K-1)*NWREC+1	0J47
	NQ=(Q-1)*NWREC+1	0J48
	DO 170 N=1,NWREC	0J49
	X =A(NK)	0J50
	A(NK)=A(NQ)	0J51
	A(NQ)=X	0J52
	NK=NK+1	0J53
	170 NQ=NQ+1	0J54
	Q=Q-1	0J55
	GO TO 2	0J56
		0J57

PAGE 02

```

16 Q=Q-1
GO TO 13
14 GO TO 3
C
L
2 CONTINUE
C
3 NI=(I-1)*NWREC+1
NQ=(Q-1)*NWREC+1
DO 300 N=1,NWREC
A(NI)=A(NQ)
A(NQ)=T(N)
NI=NI+1
300 NQ=NQ+1
IF((2*Q)-(I+J))18,18,19
19 LT(M)=I
UT(M)=Q-1
I=Q+1
GO TO 20
18 LT(M)=Q+1.
UT(M)=J
J=Q-1
20 M=M+1
GO TO 1
10 IF(I-J)21,4,4
C*****INSERT YOUR OWN CODE FOR NON NUMERIC KEYS
C*****OR KEYS WHICH ARE NOT THE INITIAL WORDS
C*****OF THE RECORD -FROM HERE
21 NI=(I-1)*NWREC+1
NJ=(J-1)*NWREC+1
DO 211 N=1,NWREC
IF (A(NI)-A(NJ))4,210,22
210 NI=NI+1
211 NJ=NJ+1
GO TO 4
C***** -TO HERE
22 NI=(I-1)*NWREC+1
NJ=(J-1)*NWREC+1
DO 220 N=1,NWREC
X=A(NI)
A(NI)=A(NJ)
A(NJ)=X
NI=NI+1
220 NJ=NJ+1
C
4 M=M+1
IF(M)24,24,23
23 I=LT(M)
J=UT(M)
GO TO 1
24 RETURN
END

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VARIABLE ALLOCATIONS  
UT =0010 LT =001F Q =0020 P =0021 X =0022 S =0023 I =0024 J =0025 NP =0026 .I =0027  
K =0028 .11 =0029 K =002A NK =002B NQ =002C NJ =002D

STATEMENT ALLOCATIONS  
1 =0072 11 =007A 110 =008E 130 =00FF 13 =010F 15 =0115 150 =013A 17 =014A 170 =0154 16 =016A  
14 =01A2 2 =01A4 3 =01AC 300 =01E7 19 =0206 18 =021F 20 =0236 10 =023E 21 =0244 210 =0274  
211 =027A 22 =028A 220 =02C4 4 =02D2 23 =02DC 24 =02ED

FEATURES SUPPORTED  
ONE WORD INTEGERS  
CALLED SUBPROGRAMS  
SUMSC SUBIN  
INTEGER CONSTANTS  
1=0032 2=0033

CORE REQUIREMENTS FOR USORT  
COMMON 0 VARIABLES 50 PROGRAM 702

END OF COMPILATION

// DUP  
\*STORE    US UA USORT  
1A92 002A

0109  
0110

// FOR

0001

\*    NAME DSRT1  
\*    ONE WORD INTEGERS  
\*    LIST ALL

PAGE 01  
0002  
0003  
0003A



```

SUBROUTINE DSRT1(NFIL,NREC,NKEY,ARRAY,TEMP)
C A SUBROUTINE WHICH SORTS (IN ASCENDING
C SEQUENCE) THE FIRST NRFILE RECORDS
C OF NREC WORDS EACH IN FILE 1.
C ARRAY IS A WORKING STORAGE VECTOR OF 2560 WORDS
C TEMP IS A WORKING STORAGE VECTOR OF NREC WORDS
C THERE ARE NKEY WORDS IN THE KEY WHICH
C ARE INTEGERS, THE FIRST IN THE RECORD
C AND IN MAJOR TO MINOR SEQUENCE
C INTEGER FLAG1,FLAG2,ARRAY(2560),TEMP(320)
NRSEC = 320 / NREC
NRH = NRSEC * 4
NRH = NRSEC * 6
C I COUNTS RECORDS IN THE FILE
C J COUNTS RECORDS IN A BATCH
C K COUNTS WORDS IN A BATCH
C L COUNTS WORDS IN A RECORD
C M IS A CONTROL INDEX
C NLH POINTS TO RECORD IN LOWER HALF
C NUH POINTS TO RECORD IN UPPER HALF
C KEYL POINTS TO A KEY WORD IN NLH'N RECORD
C KEYU POINTS TO A KEY WORD IN NUH'N RECORD
C N COUNTS WORDS IN KEY
C PHASE 1 TRANSFORM FILE INTO ORDERED
C STRINGS OF NRH RECORDS
I = 1
M = 1
FLAG1 = 0
C MORE AFTER THIS BATCH+
10 IF ((I-1)*NRH)-NRFILE) 12,11,11
11 NREC = NRFILE - (I-1)
FLAG1 = 1
GO TO 13
12 NREC = NRH
13 IT = I
C DREAD
GO TO 100
14 CALL QSORT (NREC,NWREC,NWKEY,ARRAY,TEMP)
I = IT
C DWRITE
GO TO 110
15 CONTINUE
IF (FLAG1) 99,10,19
C PHASE 1 COMPLETE
C PHASE 2 SUCCESSIVE MERGE PASSES
19 I=NRH+1
FLAG1=0
FLAG2=0
M=2
C IS THIS A FULL BATCH +
20 IF((I-1)*NRH)-NRFILE) 22,21,21
21 NREC=NRFILE-(I-1)
FLAG1=1
NRUMB=NRFILE-(I-1+NRHB)
C ONLY ONE STRING +
IF (NRUMB) 26,26,23
22 NREC = NRH
NRUMB=NRHB
23 IT=I
C DREAD
GO TO 100
24 I = I+1
C ***** INSERT YOUR OWN CODE FOR NON-NUMERIC KEYS
C ***** OR KEYS WHICH ARE NOT THE INITIAL WORDS
C ***** OF THE RECORDS ***** FROM HERE
KEYL = (NRHB-1)*NRSEC+1
KEYU = NRH*NRSEC+1
DO 25 A=1,NKEY
IF (ARRAY(KEYL+ARRHAY(KEYU))) 30,29,120
29 KEYL=KEYL+1
25 KEYU=KEYU+1
C ***** -TO HERE
30 I=1+NRHB+NRUMB
26 IF (FLAG1) 99,20,27
27 IF (FLAG2) 99,60,28
28 GO TO (99,31,19) I,M
31 I=1
FLAG1=0
FLAG2=0
M=3
GO TO 20
60 RETURN
99 PAUSE 1
CALL EXIT
C DREAD
C ROUTINE TO READ FROM DISK STARTING AT
C RECORD I NREC RECORDS OF NREC
C WORDS EACH INTO ARRAY, RETURN
C INDEX IS M
100 K=1
K'=NREC
DO 101 J=1,NREC
READ (1:1) (ARRAY(L),L=K,KM)
I=I+1
K=K+NREC
101 K'=K'+NREC
GO TO (14,24,24) I,M
C DWRITE
C ROUTINE TO WRITE ON DISK STARTING
C AT RECORD I NREC RECORDS OF NREC
C WORDS EACH INTO ARRAY, RETURN
C INDEX IS M
110 K=1
K'=NREC
DO 111 J=1,NREC
WRITE (1:1) (ARRAY(L),L=K,KM)
I=I+1
K=K+NREC
111 K'=K'+NREC
GO TO 15
C MERGE
C ROUTINE TO MERGE (NRHB RECORDS
C FROM THE LOWER HALF OF ARRAY
C WITH NRUMB RECORDS FROM THE
C UPPER HALF OF ARRAY AND WRITE
C THESE ON DISK STARTING AT
C RECORD I, NREC WORDS PER RECORD,
C NKEY WORDS IN KEY, M IS RETURN INDEX
120 NLH=1

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```

NUM=NRHB*NWREC+1
FLAG2=1
C COMPARE KEYS IN TWO RECORDS
C****INSERT YOUR OWN CODE FOR NON NUMERIC KEYS
C****OR KEYS WHICH ARE NOT THE INITIAL WORDS
C****OF THE RECORD =FROM HERE
121 KEYL=NLH
KEYU=NUM
DO 123 N=1,NWKEY
IF(ARRAY(KEYL)-ARRAY(KEYU))124,122,126
122 KEYL = KEYL+1
123 KEYU = KEYU+1
C**** -TO HERE
C RECORD FROM LOWER HALF IS WRITTEN
124 NLHM=NLH+NWREC-1
WRITE (1:1) (ARRAY(L),L=NLH,NLHM)
I=I+1
NLH=NLH+NWREC
C ARE RECORDS IN LOWER HALF EXHAUSTED
IF (NLHM/NWREC)-NRHB)121,128,99
C RECORD FROM UPPER HALF IS WRITTEN
126 NUHM = NUM+NWREC-1
WRITE (1:1) (ARRAY(L),L=NUM,NUHM)
I=I+1
NUM=NUM+NWREC
C ARE RECORDS IN UPPER HALF EXHAUSTED
IF ((NUM/NWREC)-(NRUMB+NRHB) )
1121,129,99
C RUN OUT RECORDS IN UPPER HALF
128 NLIM = (NRUMB+NRHB)*NWREC
NRO=NUM
GO TO 130
C RUN OUT RECORDS IN LOWER HALF
129 NLIM = NRHB*NWREC
NRO=NLH
C RUN OUT
130 NROM = NRO+NWREC-1
WRITE (1:1) (ARRAY(L),L=NRO,NROM)
I=I+1
NRO=NRO+NWREC
IF(NROM-NLIM) 130,26,99
C MERGE COMPLETE
END

```

PAGL 03

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0125  
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PAGE 04

VARIABLE ALLOCATIONS  
FLAG1=0002 FLAG2=0003 NRSEC=0004 NRHB =0005 NKB =0006 I =0007 M =0008 NREC =0009 IT =000A NRUMB=000B  
KEYL =000C KEYU =000D N =000E K =000F KM =0010 J =0011 L =0012 NLN =0013 NUH =0014 NLHM =0015  
NUHM =0016 NLIN =0017 NRO =0018 NROM =0019

STATEMENT ALLOCATIONS  
10 =0074 11 =007E 12 =008E 13 =0092 14 =0098 15 =00A5 19 =00AB 20 =00BD 21 =00C7 22 =00E7  
23 =00EF 24 =00F5 29 =0127 25 =012D 30 =013B 26 =0143 27 =0149 28 =014F 31 =0156 60 =0168  
90 =016A 100 =016D 101 =019B 110 =018D 111 =01DF 120 =01EF 121 =0200 122 =0222 123 =0228 124 =0236  
126 =026E 128 =02AE 129 =02BD 130 =02C8

FEATURES SUPPORTED  
ONE WORD INTEGERS

CALLED SUBPROGRAMS  
USORT SUBSC PAUSE SUBIN SDRED SDWRT SDCOM SDIX

INTEGER CONSTANTS  
320=001C 4=001D 8=001E 1=001F 0=0020 2=0021 3=0022

CORE REQUIREMENTS FOR DSHT:  
COMMON 0 VARIABLES 28 PROGRAM 738

END OF COMPILATION

// DUP

0167

\*STORE WS UA DSRT1

0168

LABC 002C

// FOR

0001

\* NAME MAINL  
\* ONE WORD INTEGERS  
\* IOCS (CARD,1132PRINTER,DISK)  
\* LIST ALL

PAGE 01

0002  
0003  
0004  
0005

```

      INTEGER ARRAY (2560),TEMP(39)
      DEFINE FILE 1 (500,39,U,IPT)
C     FIRST CARD TELLS HOW MANY CARDS FOLLOW
      READ (2,90) NRFIL
      WRITE (3,92)
C     LOOP READS CARD,WRITES RECORD ON DISK AND PRINTS RECORD
      DO 1 I=1,NRFIL
      READ (2,90) N1,N2,(TEMP(J),J=1,37)
      WRITE (1,1) N1,N2,(TEMP(J),J=1,37)
C     1 WRITE (3,91)N1,N2,(TEMP(J),J=1,37)
      SORT
      NRKEY=2
      NRREC=39
      CALL DSRT1 (NRFIL,NRREC,NRKEY,ARRAY,TEMP)
C     PRINT SORTED FILE
      WRITE (3,93)
      DO 2 I=1,NRFIL
      READ (1,1) N1,N2,(TEMP(J),J=1,37)
      2 WRITE (3,91) N1,N2,(TEMP(J),J=1,37)
      CALL EXIT
      90 FORMAT (2I3,37A2)
      91 FORMAT (1H,13,1X,13,1X,37A2)
      92 FORMAT (48H1RECORDS PRINTED AS READ
      93 FORMAT (48H0RECORDS PRINTED AS SORTED
      END
  
```

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VARIABLE ALLOCATIONS
ARRAY=0A07 TEMP =0A2E IPT =0A2F NRFIL=0A30 I =0A31 N1 =0A32 N2 =0A33 J =0A34 NRKEY=0A35 NRREC=0A36

STATEMENT ALLOCATIONS
90 =0A3D 91 =0A42 92 =0A4B 93 =0A65 1 =0A06 2 =0B2A

FEATURES SUPPORTED
ONE WORD INTEGERS
IOCS

CALLED SUBPROGRAMS
DSRT1 FLD FSTO SRED SWRT SCOMP SFIO SIOIX SIOI SUBSC CARDZ PRNTZ SDFIG SDFIX

INTEGER CONSTANTS
2=0A38 3=0A39 1=0A3A 37=0A3B 39=0A3C

CORE REQUIREMENTS FOR MAINL
COMMON 0 VARIABLES 2616 PROGRAM 278

END OF COMPILATION
  
```

// XEO

0040

RECORDS PRINTED AS HEAD  
 2 2 SORT ROUTINES WOULD BE IF THE RESTRICTIONS ARE OBSERVED. 0042  
 2 1 FOR UPDATING OR REPORT PREPARATION. THEY CAN HOWEVE BE USED AS ANY 0043  
 1 2 SO THAT ONE CAN FIND RECORDS IN A LARGER FILE IN PROPER SEQUENCE 0044  
 1 1 THESE SORTS ARE INTENDED FOR SORTING KEYS AND TAGS RECORD NUMBERS 0045

RECORDS PRINTED AS SORTED  
 1 1 THESE SORTS ARE INTENDED FOR SORTING KEYS AND TAGS RECORD NUMBERS 0045  
 1 2 SO THAT ONE CAN FIND RECORDS IN A LARGER FILE IN PROPER SEQUENCE 0046  
 2 1 FOR UPDATING OR REPORT PREPARATION. THEY CAN HOWEVE BE USED AS ANY 0043  
 2 2 SORT ROUTINES WOULD BE IF THE RESTRICTIONS ARE OBSERVED. 0042

*DSRT2, MAIN, complete solut & run*

// JOB T  
// FOR

0001

NAME DSRT2  
ONE WORD INTEGERS  
LIST ALL

PAGE 01

0002  
0003  
0003A



VARIABLE ALLOCATIONS

J	=0802	P	=0003	UT	=0013	LT	=0023	AI	=0163	AJ	=02A3	AQ	=03E3	AK	=0523	T	=0663	M	=0664
I	=0665	J	=0666	L	=0667	11	=0668	K	=0669	N	=066A								

STATEMENT ALLOCATIONS

1	=06AE	11	=0686	130	=0737	13	=0741	15	=0747	150	=0758	17	=0762	170	=0794	16	=07C3	14	=07E1
2	=07E3	3	=07EB	300	=0826	19	=0845	18	=0874	20	=0888	10	=0893	21	=0899	210	=08C0	22	=08CA
4	=08FB	23	=0902	24	=0929														

FEATURES SUPPORTED  
ONE WORD INTEGERS

CALLED SUBPROGRAMS

SUBSC	SUBIN	SDRED	SDWRT	SDCON	SDIX
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INTEGER CONSTANTS

1=066E	2=066F
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CORE REQUIREMENTS FOR DSRT2

COMMON	0	VARIABLES	1646	PROGRAM	702
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END OF COMPILATION

// DUP

\*STORE WS UA DSRT2

1A92 002D

0093

0094

// FOR

0001

\* NAME MAINL  
 \* ONE WORD INTEGERS  
 \* IOCS (CARD,1132PRINTER,DISK)  
 \* LIST ALL

PAGE 01

0002  
 0003  
 0004  
 0005

	INTEGER	TEMP(39)		PAGE 02
	DEFINE FILE 1	(500,39,U,1PT)		0006
C	FIRST CARD TELLS HOW MANY CARDS FOLLOW			0007
	READ (2,90)	NRFIL		0008
	WRITE (3,92)			0009
C	LOOP READS CARD,WRITES RECORD ON DISK AND PRINTS RECORD			0010
	DO 1 I=1,NRFIL			0011
	READ (2,90)	N1,N2,(TEMP(J),J=1,37)		0012
	WRITE (1,1)	N1,N2,(TEMP(J),J=1,37)		0013
1	WRITE (3,91)	N1,N2,(TEMP(J),J=1,37)		0014
C	SORT			0015
	NWKEY=2			0016
	NWRFC=39			0017
	CALL DSRT2	(NRFIL,NWREC,NWKEY)		0018
C	PRINT SORTED FILE			0019
	WRITE (3,93)			0020
	DO 2 I=1,NRFIL			0021
	READ (1,1)	N1,N2,(TEMP(J),J=1,37)		0022
2	WRITE (3,91)	N1,N2,(TEMP(J),J=1,37)		0023
	CALL EXIT			0024
90	FORMAT (2I3,37A2)			0025
91	FORMAT (1H,13,1X,13,1X,37A2)			0026
92	FORMAT (4RH1RECORDS PRINTED AS READ		)	0027
93	FORMAT (4RH0RECORDS PRINTED AS SORTED		)	0028
	END			0029
				0030



VARIABLE ALLOCATIONS

TEMP =002E IPT =002F NHFIL=0030 I =0031 N1 =0032 N2 =0033 J =0034 MWKEY=0035 NHREC=0036

STATEMENT ALLOCATIONS

90 =003D 91 =0042 92 =004B 93 =0065 1 =00D6 2 =012B

FEATURES SUPPORTED  
ONE WORD INTEGERS  
IOCS

CALLED SUBPROGRAMS

DSRTZ FLD FSTO SRED SWRT SCOMP SF10 S101X S101 SUBSC CARDZ PRNTZ SDFIO SDRED SWRT  
SDCOM SDIX SDI

INTEGER CONSTANTS

2=0038 3=0039 1=003A 37=003B 39=003C

CORE REQUIREMENTS FOR MAIN

COMMON 0 VARIABLES 56 PROGRAM 276

END OF COMPILATION

// XEQ

0001

RECORDS PRINTED AS READ

2 2 SORT ROUTINES WOULD BE IF THE RESTRICTIONS ARE OBSERVED. 0003  
2 1 FOR UPDATING OR REPORT PREPARATION. THEY CAN HOWEVER BE USED AS ANY 0004  
1 2 SO THAT ONE CAN FIND RECORDS IN A LARGER FILE IN PROPER SEQUENCE 0005  
1 1 THESE SORTS ARE INTENDED FOR SORTING KEYS AND TAGS RECORD NUMBERS 0006

RECORDS PRINTED AS SORTED

1 1 THESE SORTS ARE INTENDED FOR SORTING KEYS AND TAGS RECORD NUMBERS 0006  
1 2 SO THAT ONE CAN FIND RECORDS IN A LARGER FILE IN PROPER SEQUENCE 0005  
2 1 FOR UPDATING OR REPORT PREPARATION. THEY CAN HOWEVER BE USED AS ANY 0004  
2 2 SORT ROUTINES WOULD BE IF THE RESTRICTIONS ARE OBSERVED. 0003