

FILE ID**VAXARITH

c 7

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:LJK0045
-1

0000 1 .TITLE VAX\$DECIMAL_ARITHMETIC - VAX-11 Packed Decimal Arithmetic Instructio
0000 .1 .IDENT /V04-001/
0000 3
0000 4 :
0000 5 :*****
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0000 24 :**
0000 25 :**
0000 26 :*****
0000 27 :
0000 28 :
0000 29 :++
0000 30 : Facility:
0000 31 :
0000 32 : VAX-11 Instruction Emulator
0000 33 :
0000 34 : Abstract:
0000 35 :
0000 36 : The routines in this module emulate the VAX-11 packed decimal
0000 37 : instructions that perform arithmetic operations. These procedures can
0000 38 : be a part of an emulator package or can be called directly after the
0000 39 : input parameters have been loaded into the architectural registers.
0000 40 :
0000 41 : The input parameters to these routines are the registers that
0000 42 : contain the intermediate instruction state.
0000 43 :
0000 44 : Environment:
0000 45 :
0000 46 : These routines run at any access mode, at any IPL, and are AST
0000 47 : reentrant.
0000 48 :
0000 49 : Author:
0000 50 :
0000 51 : Lawrence J. Kenah
0000 52 :
0000 53 : Creation Date
0000 54 :
0000 55 : 19 October 1983
0000 56 :
0000 57 : Modified by:

:LJK0045	0000	58	:	
:LJK0045	0000	.1	:	V04-001 LJK0045 Lawrence J. Kenah 19-Sep-1984
:LJK0045	0000	.2	:	The result string in ADDP4 and SUBP4 must be probed for
:LJK0045	0000	.3	:	write access to insure that the instruction does not modify
:LJK0045	0000	.4	:	memory before trying to restart.
:LJK0045	0000	.5	:	
	0000	59	:	V01-003 LJK0037 Lawrence J. Kenah 17-Jul-1984
	0000	60	:	Fix two minor bugs in exception handling code that caused
	0000	61	:	MULP and DIVP tests to generate spurious access violations.
	0000	62	:	
	0000	63	:	V01-002 LJK0024 Lawrence J. Kenah 21-Feb-1984
	0000	64	:	Add code to handle access violations. Perform minor cleanup.
	0000	65	:	Eliminate double use of R10 in MULP and DIVP.
	0000	66	:	
	0000	67	:	V01-001 LJK0008 Lawrence J. Kenah 19-Oct-1983
	0000	68	:	The emulation code for ADDP4, ADDP6, SUBP4, SUBP6, MULP and
	0000	69	:	DIVP was moved into a separate module.
	0000	70	--	

```

0000 72 .SUBTITLE Declarations
0000 73
0000 74 : Include files:
0000 75
0000 76 .NOCROSS
0000 77 .ENABLE SUPPRESSION ; No cross reference for these
0000 78 ; No symbol table entries either
0000 79 ADDP4_DEF ; Bit fields in ADDP4 registers
0000 80 ADDP6_DEF ; Bit fields in ADDP6 registers
0000 81 DIVP_DEF ; Bit fields in DIVP registers
0000 82 MULP_DEF ; Bit fields in MULP registers
0000 83 SUBP4_DEF ; Bit fields in SUBP4 registers
0000 84 SUBP6_DEF ; Bit fields in SUBP6 registers
0000 85
0000 86 $PSLDEF ; Define bit fields in PSL
0000 87 $SRMDEF ; Define arithmetic trap codes
0000 88
0000 89 .DISABLE SUPPRESSION ; Turn on symbol table again
0000 90 .CROSS ; Cross reference is OK now
0000 91
0000 92 : Symbol definitions
0000 93
0000 94 : The architecture requires that R4 be zero on completion of an ADDP6 or
0000 95 : SUBP6 instruction. If we did not have to worry about restarting
0000 96 : instructions after an access violation, we could simply zero the saved
0000 97 : R4 value on the code path that these two instructions have in common
0000 98 : before they merge with the ADDP4 and SUBP4 routines. The ability to
0000 99 : restart requires that we keep the original R4 around at least until no
0000 100 : more access violations are possible. To accomplish this, we store the
0000 101 : fact that R4 must be cleared on exit in R11, which also contains the
0000 102 : evolving condition codes. We use bit 31, the compatibility mode bit
0000 103 : because it is nearly impossible to enter the emulator with CM set.
0000 104
0000 0000001F 105 ADD_SUB_V_ZERO_R4 = PSL$V_CM
0000 106
0000 107 : External declarations
0000 108
0000 109 .DISABLE GLOBAL
0000 110
0000 111 .EXTERNAL -
0000 112 DECIMAL$BOUNDS_CHECK,-
0000 113 DECIMAL$BINARY_TO_PACKED_TABLE,-
0000 114 DECIMAL$PACKED_TO_BINARY_TABLE,-
0000 115 DECIMAL$STRIP_ZEROS_R0_RT,-
0000 116 DECIMAL$STRIP_ZEROS_R2_R3
0000 117
0000 118 .EXTERNAL -
0000 119 VAX$DECIMAL_EXIT,-
0000 120 VAX$DECIMAL_ACCVIO,-
0000 121 VAX$REFLECT_TRAP,-
0000 122 VAX$ROPRAND
0000 123
0000 124 : PSECT Declarations:
0000 125
0000 126 .DEFAULT DISPLACEMENT , WORD
0000 127
0000 128 .PSECT _VAX$CODE PIC, USR, CON, REL, LCL, SHR, EXE, RD, NOWRT, LONG

```

VAX\$DECIMAL_ARITHMETIC
V04-001

- VAX-11 Packed Decimal Arithmetic Instr ^{H 7}
Declarations 8-JAN-1985 17:27:01 VAX/VMS Macro V04-00
5-SEP-1984 00:44:34 [EMULAT.BUGSRC]VAXARITH.MAR;1 Page 4
(2)

0000 129
0000 130

BEGIN_MARK_POINT

-1

```

0000 133 .SUBTITLE      VAX$SUBP6 - Subtract Packed (6 Operand Format)
0000 134 :+
0000 135 : Functional Description:
0000 136 :
0000 137 : In 6 operand format, the subtrahend string specified by the subtrahend
0000 138 : length and subtrahend address operands is subtracted from the minuend
0000 139 : string specified by the minuend length and minuend address operands.
0000 140 : The difference string specified by the difference length and difference
0000 141 : address operands is replaced by the result.
0000 142 :
0000 143 : Input Parameters:
0000 144 :
0000 145 : R0 - sublen.rw      Number of digits in subtrahend string
0000 146 : R1 - subaddr.ab    Address of subtrahend string
0000 147 : R2 - minlen.rw    Number of digits in minuend string
0000 148 : R3 - minaddr.ab  Address of minuend string
0000 149 : R4 - diflen.rw    Number of digits in difference string
0000 150 : R5 - difaddr.ab  Address of difference string
0000 151 :
0000 152 : Output Parameters:
0000 153 :
0000 154 : R0 = 0
0000 155 : R1 = Address of the byte containing the most significant digit of
0000 156 :   the subtrahend string
0000 157 : R2 = 0
0000 158 : R3 = Address of the byte containing the most significant digit of
0000 159 :   the minuend string
0000 160 : R4 = 0
0000 161 : R5 = Address of the byte containing the most significant digit of
0000 162 :   the string containing the difference
0000 163 :
0000 164 : Condition Codes:
0000 165 :
0000 166 : N <- difference string LSS 0
0000 167 : Z <- difference string EQL 0
0000 168 : V <- decimal overflow
0000 169 : C <- 0
0000 170 :
0000 171 : Register Usage:
0000 172 :
0000 173 : This routine uses all of the general registers. The condition codes
0000 174 : are recorded in R11 as the routine executes.
0000 175 :-
0000 176 :
0000 177 .ENABLE      LOCAL_BLOCK
0000 178 :
0000 179 VAX$SUBP6:::
0000 180 PUSHR  #^M<R0,R1,R2,R3,R4,R5,R6,R7,R8,R9,R10,R11> ; Save the lot
0000 181 MOVZBL #1,R9          ; Indicate that this is subtraction
0000 182 BRB    10$            ; Merge with ADDP6 code

```

0FFF	8F	BB	0000
59	01	9A	0004
06	11		0007

0009 184 .SUBTITLE VAXSADDP6 - Add Packed (6 Operand Format)
 0009 185 :+
 0009 186 : Functional Description:
 0009 187 :
 0009 188 : In 6 operand format, the addend 1 string specified by the addend 1
 0009 189 : length and addend 1 address operands is added to the addend 2 string
 0009 190 : specified by the addend 2 length and addend 2 address operands. The sum
 0009 191 : string specified by the sum length and sum address operands is replaced
 0009 192 : by the result.
 0009 193 :
 0009 194 : Input Parameters:
 0009 195 :
 0009 196 : R0 - add1len.rw Number of digits in first addend string
 0009 197 : R1 - add1addr.ab Address of first addend string
 0009 198 : R2 - add2len.rw Number of digits in second addend string
 0009 199 : R3 - add2addr.ab Address of second addend string
 0009 200 : R4 - sumlen.rw Number of digits in sum string
 0009 201 : R5 - sumaddr.ab Address of sum string
 0009 202 :
 0009 203 : Output Parameters:
 0009 204 :
 0009 205 : R0 = 0
 0009 206 : R1 = Address of the byte containing the most significant digit of
 0009 207 : the first addend string
 0009 208 : R2 = 0
 0009 209 : R3 = Address of the byte containing the most significant digit of
 0009 210 : the second addend string
 0009 211 : R4 = 0
 0009 212 : R5 = Address of the byte containing the most significant digit of
 0009 213 : the string containing the sum
 0009 214 :
 0009 215 : Condition Codes:
 0009 216 :
 0009 217 : N <- sum string LSS 0
 0009 218 : Z <- sum string EQL 0
 0009 219 : V <- decimal overflow
 0009 220 : C <- 0
 0009 221 :
 0009 222 : Register Usage:
 0009 223 :
 0009 224 : This routine uses all of the general registers. The condition codes
 0009 225 : are recorded in R11 as the routine executes.
 0009 226 :-
 0009 227 :
 0009 228 : VAXSADDP6::
 0FFF 8F BB 0009 229 : PUSHR #^M<R0,R1,R2,R3,R4,R5,R6,R7,R8,R9,R10,R11> ; Save the lot
 59 D4 000D 230 : CLRL R9 ; This is addition
 5B DC 001A 231 10\$: ROPRAND_CHECK R4 ; Insure that R4 is !EQU 31
 001C 232 : MOVPSL R11 ; Get initial PSL
 001C 233 :
 001C 234 : Indicate that the saved R4 must be cleared on the exit path
 001C 235 :
 22 5B 1F E3 001C .1 BBCS #ADD_SUB_V_ZERO_R4,R11,25\$; Set bit and join common code
 20 11 0020 .2 BRB 25\$; In case we drop through BBCS

:LJK0045
:LJK0045

```

0022 239 .SUBTITLE      VAX$SUBP4 - Subtract Packed (4 Operand Format)
0022 240 :+
0022 241 : Functional Description:
0022 242 :
0022 243 : In 4 operand format, the subtrahend string specified by subtrahend
0022 244 : length and subtrahend address operands is subtracted from the difference
0022 245 : string specified by the difference length and difference address
0022 246 : operands and the difference string is replaced by the result.
0022 247 :
0022 248 : Input Parameters:
0022 249 :
0022 250 R0 - sublen.rw      Number of digits in subtrahend string
0022 251 R1 - subaddr.ab    Address of subtrahend decimal string
0022 252 R2 - diflen.rw      Number of digits in difference string
0022 253 R3 - difaddr.ab    Address of difference decimal string
0022 254 :
0022 255 : Output Parameters:
0022 256 :
0022 257 R0 = 0
0022 258 R1 = Address of the byte containing the most significant digit of
0022 259 : the subtrahend string
0022 260 R2 = 0
0022 261 R3 = Address of the byte containing the most significant digit of
0022 262 : the string containing the difference
0022 263 :
0022 264 : Condition Codes:
0022 265 :
0022 266 N <- difference string LSS 0
0022 267 Z <- difference string EQL 0
0022 268 V <- decimal overflow
0022 269 C <- 0
0022 270 :
0022 271 : Register Usage:
0022 272 :
0022 273 This routine uses all of the general registers. The condition codes
0022 274 : are recorded in R11 as the routine executes.
0022 275 :-
0022 276 :
0022 277 VAX$SUBP4::          PUSHR #^M<R0,R1,R2,R3,R4,R5,R6,R7,R8,R9,R10,R11> ; Save the lot
1FFF 8F BB 0022 278          MOVZBL #1,R9 ; Indicate that this is subtraction
59 01 9A 0026 279          BRB 20$ ; Merge with ADDP4 code
0B 11 0029 280

```

002B 282 .SUBTITLE VAX\$ADDP4 - Add Packed (4 Operand Format)
 002B 283 :+
 002B 284 : Functional Description:
 002B 285 :
 002B 286 : In 4 operand format, the addend string specified by the addend length
 002B 287 : and addend address operands is added to the sum string specified by the
 002B 288 : sum length and sum address operands and the sum string is replaced by
 002B 289 : the result.
 002B 290 :
 002B 291 : Input Parameters:
 002B 292 :
 002B 293 : R0 - addlen.rw Number of digits in addend string
 002B 294 : R1 - addaddr.ab Address of addend decimal string
 002B 295 : R2 - sumlen.rw Number of digits in sum string
 002B 296 : R3 - sumaddr.ab Address of sum decimal string
 002B 297 :
 002B 298 : Output Parameters:
 002B 299 :
 002B 300 : R0 = 0
 002B 301 : R1 = Address of the byte containing the most significant digit of
 002B 302 : the addend string
 002B 303 : R2 = 0
 002B 304 : R3 = Address of the byte containing the most significant digit of
 002B 305 : the string containing the sum
 002B 306 :
 002B 307 : Condition Codes:
 002B 308 :
 002B 309 : N <- sum string LSS 0
 002B 310 : Z <- sum string EQL 0
 002B 311 : V <- decimal overflow
 002B 312 : C <- 0
 002B 313 :
 002B 314 : Register Usage:
 002B 315 :
 002B 316 : This routine uses all of the general registers. The condition codes
 002B 317 : are recorded in R11 as the routine executes.
 002B 318 :-
 002B 319 :
 ;LJK0045 0264 30 002B .1 15\$: BSBW CHECK_WRITE_ACCESS : Perform rigorous access check
 ;LJK0045 38 11 002E .2 BRB 30\$: String can be written after all
 ;LJK0045 0030 0030 .3
 0FFF BF 88 0030 320 VAX\$ADDP4:::
 59 D4 0034 321 PUSHR #^M<R0,R1,R2,R3,R4,R5,R6,R7,R8,R9,R10,R11> : Save the lot
 0034 322 CLRL R9 ; This is addition
 0036 323
 0036 324 : The output string, described by R4 and R5, will be the same as the input
 0036 325 : string for ADDP4 and SUBP4. It is necessary to explicitly clear R4<31:16>
 0036 326 : along this code path so MOVA R2,R4 will not always work.
 0036 327
 54 52 3C 0036 328 20\$: MOVZWL R2,R4 : Set output size equal to input size
 55 53 D0 0039 329 MOVL R3,R5 : ... and ditto for string addresses
 58 DC 003C 330 MOVPSL R11 ; Get initial PSL
 003E 331
 003E 332 : Indicate that the saved R4 will be restored on the common exit path
 003E 333
 ;LJK0045 00 5B 1F E5 003E .1 BBCC #ADD_SUB_V_ZERO_R4,R11,25\$: Clear bit and join common code

0042 336 .SUBTITLE ADDPx/SUBPx Common Initialization Code
 0042 337 :+
 0042 338 : All four routines converge at this point and execute common initialization
 0042 339 : code until a later decision is made to do addition or subtraction.
 0042 340 :
 0042 341 : R4 - Number of digits in destination string
 0042 342 : R5 - Address of destination string
 0042 343 :
 0042 344 : R9 - Indicates whether operation is addition or subtraction
 0042 345 : 0 => addition
 0042 346 : 1 => subtraction
 0042 347 :
 0042 348 : R11<31> - Indicates whether this is a 4-operand or 6-operand instruction
 0042 349 : 0 => 4-operand (restore saved R4 on exit)
 0042 350 : 1 => 6-operand (set R4 to zero on exit)
 0042 351 :-
 0042 352 :
 ;LJK0045 04 00 04 F0 0042 .1 258: INSV #PSLSM_Z,#0,#4,R11 ; Set Z-bit, clear the rest in saved PSW
 -1 5B 0046 0047 354 ESTABLISH_HANDLER - ; Store address of access
 0047 355 ARITH_ACCVIO ; violation handler
 004C 356 :
 004C 357 ROPRAND CHECK R2 : Insure that R2 is LEQU 31
 0054 358 MARK_POINT ADD SUB BSBW 0
 FFA9' 30 0054 359 BSBW DECIMAL\$STRTP_ZEROS_R2_R3 : Strip high order zeros from R2/R3
 0057 360 :
 0057 361 ROPRAND CHECK R0 : Insure that R0 is LEQU 31
 005F 362 MARK_POINT ADD SUB BSBW 0
 FF9E' 30 005F 363 BSBW DECIMAL\$STRIP_ZEROS_R0_R1 : Strip high order zeros from R0/R1
 0062 364 :
 ;LJK0045 0062 .1 : Perform the access check on the output string for the worst case, a string
 ;LJK0045 0062 .2 : large enough to accommodate 31 decimal digits. A detailed check using the
 ;LJK0045 0062 .3 : correct byte length of the output string is necessary only if this initial
 ;LJK0045 0062 .4 : probe fails. The more detailed check can be handled out of line.
 ;LJK0045 0062 .5 :
 ;LJK0045 65 10 00 0D 0062 .6 PROBEW #0,#16,(R5) ; Can result string be written?
 ;LJK0045 C3 13 0066 .7 BEQL 15\$; Branch if no write access allowed
 ;LJK0045 0068 0068 365 : Rather than totally confuse the already complicated logic dealing with
 ;LJK0045 0068 366 : different length strings in the add or subtract loop, we will put the
 ;LJK0045 0068 367 : result into an intermediate buffer on the stack. This buffer will be long
 ;LJK0045 0068 368 : enough to handle the worst case so that the addition loop need only concern
 ;LJK0045 0068 369 : itself with the lengths of the two input loops. The required length is 17
 ;LJK0045 0068 370 : bytes to handle an addition with a carry out of the most significant byte.
 ;LJK0045 0068 371 : We will allocate 20 bytes to maintain whatever alignment the stack has.
 ;LJK0045 0068 372 :
 ;LJK0045 7E 7C 0068 .1 30\$: CLRQ -(SP) ; Set aside space for output string
 -1 7E 7C 006A 374 CLRQ -(SP) ; Worst case string needs 16 bytes
 7E D4 006C 375 CLRL -(SP) ; Add slack for a CARRY
 54 04 01 EF 006E 376 EXTZV #1,#4,R4,R8 ; Get byte count of destination string
 58 0072 :
 7E 55 58 C1 0073 377 ADDL3 R8,R5,-(SP) ; Save high address end of destination
 55 18 AE 9E 0077 378 MOVAB 24(SP),R5 ; Point R5 one byte beyond buffer
 0078 379 :
 0078 380 : The number of minus signs will determine whether the real operation that we
 0078 381 : perform is addition or subtraction. That is, two plus signs or two minus
 0078 382 : signs will both result in addition, while a plus sign and a minus sign will

007B 383 : result in subtraction. The addition and subtraction routines have their own
 007B 384 : methods for determining the correct sign of the result.
 007B 385 :
 007B 386 : For the purpose of counting minus signs, we treat subtraction as the
 007B 387 : addition of the negative of the input operand. That is, subtraction of a
 007B 388 : positive quantity causes the sign to be remembered as minus and counted as
 007B 389 : a minus sign while subtraction of a minus quantity stores a plus sign and
 007B 390 : counts nothing.
 007B 391 :
 007B 392 : On input to this code sequence, R9 distinguished addition from subtraction.
 007B 393 : On output, it contains either 0, 1, or 2, indicating the total number of
 007B 394 : minus signs, real or implied, that we counted.
 007B 395 :

50 04 01 EF	007B 396 EXTZV #1,#4,R0,R6	; Get byte count for first input string
56 007F		
51 56 C0	0080 397 ADDL R6,R1	; Point R1 to byte containing sign
61 F0 8F 88	0083 398 MARK POINT ADD SUB 24	
56 0087	399 BICB3 #^B11110000,(R1),R6	; R6 contains the sign "digit"
10 59 E8	0088 400 BLBS R9,35\$; Use second CASE if subtraction
	0088 401	
	0088 402 ; This case statement is used for addition	
	0088 403	
	0088 404 CASE R6,TYPE=B,LIMIT=#10,<-	; Dispatch on sign digit
	405 50\$,-	; 10 => sign is '+'
	406 40\$,-	; 11 => sign is '-'
	407 50\$,-	; 12 => sign is '+'
	408 40\$,-	; 13 => sign is '-'
	409 50\$,-	; 14 => sign is '+'
	410 50\$,-	; 15 => sign is '-'
	411 >	
	0098 412	
	0098 413 ; This case statement is used for subtraction	
	0098 414	
	0098 415 35\$: CASE R6,TYPE=B,LIMIT=#10,<-	; Dispatch on sign digit
	416 40\$,-	; 10 => treat sign as '-'
	417 50\$,-	; 11 => treat sign as '+'
	418 40\$,-	; 12 => treat sign as '-'
	419 50\$,-	; 13 => treat sign as '+'
	420 40\$,-	; 14 => treat sign as '-'
	421 40\$,-	; 15 => treat sign as '+'
	422 >	
	00AB 423	
59 01 D0	00AB 424 40\$: MOVL #1,R9	; Count a minus sign
56 0D 9A	00AE 425 MOVZBL #13,R6	; The preferred minus sign is 13
05 11	00B1 426 BRB 60\$; Now check second input sign
	00B3 427	
56 59 D4	00B3 428 50\$: CLRL R9	; No real minus signs so far
0C 9A	00B5 429 MOVZBL #12,R6	; The preferred minus sign is 12
0088 430		
52 04 01 EF	00B8 431 60\$: EXTZV #1,#4,R2,R7	; Get byte count for second input string
57 00BC		
53 57 C0	00BD 432 ADDL R7,R3	; Point R3 to byte containing sign
63 F0 8F 88	00C0 433 MARK POINT ADD SUB 24	
57 00C4	434 BICB3 #^B11110000,(R3),R7	; R7 contains the sign "digit"
00C5 435		

00C5 436 CASE R7 TYPE=B,LIMIT=#10,<- ; Dispatch on sign digit
00C5 437 80\$,- ; 10 => sign is +
00C5 438 70\$,- ; 11 => sign is -
00C5 439 80\$,- ; 12 => sign is +
00C5 440 70\$,- ; 13 => sign is -
00C5 441 80\$,- ; 14 => sign is +
00C5 442 80\$,- ; 15 => sign is -
00C5 443 >
00D5 444
57 59 D6 00D5 445 70\$: INCL R9 ; Remember that sign was minus
0D 9A 00D7 446 MOVZBL #13,R7 ; The preferred minus sign is 13
03 11 00DA 447 BRB 90\$; Now check second input sign
57 0C 9A 00DC 449 80\$: MOVZBL #12,R7 ; The preferred minus sign is 12
00DF 450
03 59 E9 00DF 451 90\$: BLBC R9,ADD_PACKED ; Even parity indicates addition
00E2 452
00B3 31 00E2 453 BRW SUBTRACT_PACKED ; Odd parity calls for subtraction
00E5 454
00E5 455 .DISABLE LOCAL_BLOCK

00E5 457 .SUBTITLE ADD_PACKED - Add Two Packed Decimal Strings
 00E5 458 :+
 00E5 459 Functional Description:
 00E5 460
 00E5 461 This routine adds two packed decimal strings whose descriptors
 00E5 462 are passed as input parameters and places their sum into another
 00E5 463 (perhaps identical) packed decimal string.
 00E5 464
 00E5 465 At the present time, the result is placed into a 16-byte storage
 00E5 466 area while the sum is being evaluated. This drastically reduces
 00E5 467 the number of different cases that must be dealt with as each
 00E5 468 pair of bytes in the two input strings is added.
 00E5 469
 00E5 470 The signs of the two input strings have already been dealt with
 00E5 471 so this routine performs addition in all cases, even if the original
 00E5 472 entry was at SUBP4 or SUBP6. The cases that arrive in this routine
 00E5 473 are as follows.
 00E5 474
 00E5 475 R2/R3 R0/R1 result
 00E5 476 +-----+-----+-----+
 00E5 477 | R2/R3 + R0/R1 | plus | plus | plus |
 00E5 478 |-----+-----+-----+
 00E5 479 | R2/R3 + R0/R1 | minus | minus | minus |
 00E5 480 |-----+-----+-----+
 00E5 481 | R2/R3 - R0/R1 | minus | plus | minus |
 00E5 482 |-----+-----+-----+
 00E5 483 | R2/R3 - R0/R1 | plus | minus | plus |
 00E5 484 |-----+-----+-----+
 00E5 485 | R2/R3 - R0/R1 | plus | plus | minus |
 00E5 486 |-----+-----+-----+
 00E5 487 | R2/R3 - R0/R1 | minus | minus | plus |
 00E5 488 |-----+-----+-----+
 00E5 489 | R2/R3 - R0/R1 | plus | minus | plus |
 00E5 490 |-----+-----+-----+
 00E5 491 | R2/R3 - R0/R1 | minus | plus | minus |
 00E5 492 |-----+-----+-----+
 00E5 493 | R2/R3 - R0/R1 | plus | plus | minus |
 00E5 494 Note that the correct choice of sign in all four cases is the sign
 00E5 495 of the second input string, the one described by R2 and R3.
 00E5 496
 00E5 497 Input Parameters:
 00E5 498
 00E5 499 R0<4:0> - Number of digits in first input decimal string
 00E5 500 R1 - Address of least significant digit of first input
 00E5 501 decimal string (the byte containing the sign)
 00E5 502
 00E5 503 R2<4:0> - Number of digits in second input decimal string
 00E5 504 R3 - Address of least significant digit of second input
 00E5 505 decimal string (the byte containing the sign)
 00E5 506
 00E5 507 R4<4:0> - Number of digits in output decimal string
 00E5 508 R5 - Address of one byte beyond least significant digit of
 00E5 509 intermediate string stored on the stack
 00E5 510
 00E5 511 R6<3:0> - Sign of first input string in preferred form
 00E5 512 R7<3:0> - Sign of second input string in preferred form
 00E5 513

00E5 514 : R11 - Saved PSL (Z-bit is set, other condition codes are clear)
 00E5 515 :
 00E5 516 : (SP) - Saved R5, address of least significant digit of ultimate
 00E5 517 : destination string.
 00E5 518 : 4(SP) - Beginning of 20-byte buffer to hold intermediate result
 00E5 519 :
 00E5 520 : Output Parameters:
 00E5 521 :
 00E5 522 : The particular input operation (ADDPx or SUBPx) is completed in
 00E5 523 : this routine. See the routine headers for the four routines that
 00E5 524 : request addition or subtraction for a list of output parameters
 00E5 525 : from this routine.
 00E5 526 :-
 00E5 527 :
 00E5 528 ADD_PACKED:
 59 57 90 00E5 529 MOVB R7,R9 : Use sign of second string for output
 03 59 E9 00E8 530 BLBC R9,10\$: Check if sign is negative
 5B 08 88 00EB 531 BISB #PSLSM_N,R11 : ... so the saved N-bit can be set
 00EE 532 :
 56 61 0F 88 00EE 533 MARK_POINT ADD_SUB_24
 00F2 534 10\$: BICB3 #^B00001111,(R1),R6 : Get least significant digit to R6
 57 63 0F 88 00F2 535 MARK_POINT ADD_SUB_24
 58 D4 00F6 536 BICB3 #^B00001111,(R3),R7 : Get least significant digit to R7
 0075 30 00F8 537 CLRL R8 : Start the add with CARRY off
 00FB 538 BSBW ADD_PACKED_BYT_R6_R7 : Add the two low order digits
 00FB 539 :
 00FB 540 : The following set of instructions computes the number of bytes in the two
 00FB 541 : strings and, if necessary, performs a switch so that R0 and R1 always
 00FB 542 : describe the shorter of the two strings.
 00FB 543 :
 50 04 01 EF 00FB 544 EXTZV #1,#4,R0,R0 : Convert digit count to byte count
 50 04 01 EF 0100 545 EXTZV #1,#4,R2,R2 : Do it for both strings
 52 0104 :
 52 50 D1 0105 546 CMPL R0,R2 : We want to compare the byte counts
 09 1B 0108 547 BLEQU 20\$: Skip the swap if we're already correct
 56 50 7D 010A 548 MOVQ R0,R6 : Save the longer
 50 52 7D 010D 549 MOVQ R2,R0 : Store the shorter on R0 and R1
 52 56 7D 0110 550 MOVQ R6,R2 : ... and store the longer in R2 and R3
 52 50 C2 0113 551 20\$: SUBL R0,R2 : Make R2 a difference (R2 GEQU 0)
 0116 :
 0116 552 : R0 now contains the number of bytes remaining in the shorter string.
 0116 553 : R2 contains the difference in bytes between the two input strings.
 0116 554 :
 0116 555 :
 50 D5 0116 556 TSTL R0 : Does shorter string have any room?
 06 13 0118 557 BEQL 40\$: Skip loop if no room at all
 011A 558 :
 FA 50 30 011A 559 30\$: BSBW ADD_PACKED_BYT_STRING : Add the next two bytes together
 50 F5 011D 560 SOBGTR R0,30\$: Check for end of loop
 0120 :
 52 D5 0120 561 40\$: TSTL R2 : Does longer string have any room?
 16 13 0122 562 BEQL 70\$: Skip next loops if all done
 0124 :
 0D 58 E9 0124 563 50\$: BLBC R8,60\$: Life is simple if CARRY clear
 0127 :
 56 D4 0127 564 CLRL R6 : Otherwise, CARRY must propagate
 0129 565 MARK_POINT ADD_SUB_24

57 73 9A 0129 569 MOVZBL -(R3),R7 : So add CARRY to single string
 0041 30 012C 570 BSBW ADD_PACKED_BYTE_R6_R7 : Use the special entry point
 F2 52 F5 012F 571 SOBGTR R2,50\$: Check for this string exhausted
 06 11 0132 572 BRB 70\$: Join common completion code
 0134 573
 75 73 90 0134 574 MARK_POINT ADD_SUB_24
 FA 52 F5 0137 575 60\$: MOVB -(R3),-(R5) : Simply move src to dst if no CARRY
 013A 576 SOBGTR R2,60\$: ... until we're all done
 75 58 90 013A 577
 013D 578 70\$: MOVB R8,-(R5) : Store the final CARRY
 013D 579
 013D 580 :+
 013D 581 : At this point, the result has been computed. That result must be moved to
 013D 582 : its ultimate destination, noting whether any nonzero digits are stored
 013D 583 : so that the Z-bit will have its correct setting.
 013D 584
 013D 585 : Input Parameters:
 013D 586
 013D 587 : R9<7:0> - Sign of result in preferred form
 013D 588 : R11<3:0> - Saved condition codes
 013D 589 : R11<31> - Indicates whether to set saved R4 to zero
 013D 590
 013D 591 : (SP) - Saved R5, high address end of destination string
 013D 592 :-
 013D 593
 013D 594 ADD_SUBTRACT_EXIT:
 55 6E 01 C1 013D 595 ADDL3 #1,(SP),R5 : Point R5 beyond real destination
 51 18 AE 9E 0141 596 MOVAB 24(SP),R1 : R1 locates the saved result
 010C 30 0145 597 BSBW STORE RESULT : Store the result and record the Z-bit
 12 5B 02 E0 0148 598 BBS #PSL\$V_Z,R11,100\$: Step out of line for minus zero check
 014C 599
 014C 600
 04 00 59 F0 014C 601 80\$: MARK_POINT ADD_SUB_24
 9E 0150
 03 5E 14 C0 0151 602 ADDL #20,SP : Get rid of intermediate buffer
 5B 1F E1 0154 603 BBC #ADD_SUB_V_ZERO_R4,R11,90\$: Branch if 4-operand opcode
 10 AE D4 0158 604 CLRL 16(SP) : Clear saved R4 to return zero
 FEA2' 31 015B 605 90\$: BRW VAX\$DECIMAL_EXIT : Exit through common code path
 015E 606
 015E 607 : If the result is negative zero, then the N-bit is cleared and the sign
 015E 608 : is changed to a plus sign.
 015E 609
 E7 5B 08 8A 015E 610 100\$: BICB #PSLSM_N,R11 : Clear the N-bit unconditionally
 5B 01 E0 0161 611 BBS #PSLSV_V,R11,80\$: Do not change the sign on overflow
 59 0C 90 0165 612 MOVB #12,R9 : Make sure that the sign is plus
 E2 11 0168 613 BRB 80\$: ... and rejoin the exit code

016A 615 .SUBTITLE ADD_PACKED_BYT E - Add Two Bytes Containing Decimal Digits
016A 616 :+
016A 617 : Functional Description:
016A 618 :
016A 619 : This routine adds together two bytes containing decimal digits and
016A 620 : produces a byte containing the sum that is stored in the output
016A 621 : string. Each of the input bytes is converted to a binary number
016A 622 : (with a table-driven conversion), the two numbers are added, and
016A 623 : the sum is converted back to two decimal digits stored in a byte.
016A 624 :
016A 625 : This routine makes no provisions for bytes that contain illegal
016A 626 : decimal digits. We are using the UNPREDICTABLE statement in the
016A 627 : architectural description of the decimal instructions to its fullest.
016A 628 :
016A 629 : The bytes that contain a pair of packed decimal digits can either
016A 630 : exist in packed decimal strings located by R1 and R3 or they can
016A 631 : be stored directly in registers. In the former case, the digits must
016A 632 : be extracted from registers before they can be used in later operations
016A 633 : because the sum will be used as an index register.
016A 634 :
016A 635 : For entry at ADD_PACKED_BYT E_STRING:
016A 636 :
016A 637 : Input Parameters:
016A 638 :
016A 639 : R1 - Address one byte beyond first byte that is to be added
016A 640 : R3 - Address one byte beyond second byte that is to be added
016A 641 : R5 - Address one byte beyond location to store sum
016A 642 :
016A 643 : R8 - Carry from previous byte (R8 is either 0 or 1)
016A 644 :
016A 645 : Implicit Input:
016A 646 :
016A 647 : R6 - Scratch
016A 648 : R7 - Scratch
016A 649 :
016A 650 : Output Parameters:
016A 651 :
016A 652 : R1 - Decreased by one to point to current byte in first input string
016A 653 : R3 - Decreased by one to point to current byte in second input strin
016A 654 : R5 - Decreased by one to point to current byte in output string
016A 655 :
016A 656 : R8 - Either 0 or 1, reflecting whether this most recent ADD resulted
016A 657 : in a CARRY to the next byte.
016A 658 :
016A 659 : For entry at ADD_PACKED_BYT E_R6_R7:
016A 660 :
016A 661 : Input Parameters:
016A 662 :
016A 663 : R6 - First byte containing decimal digit pair
016A 664 : R7 - Second byte containing decimal digit pair
016A 665 :
016A 666 : R5 - Address one byte beyond location to store sum
016A 667 :
016A 668 : R8 - Carry from previous byte (R8 is either 0 or 1)
016A 669 :
016A 670 : Output Parameters:
016A 671 :

016A 672 : R5 - Decreased by one to point to current byte in output string
 016A 673 :
 016A 674 : R8 - Either 0 or 1, reflecting whether this most recent ADD resulted
 016A 675 : in a CARRY to the next byte.
 016A 676 :
 016A 677 : Side Effects:
 016A 678 : R6 and R7 are modified by this routine
 016A 680 :
 016A 681 : R0, R2, R4, and R9 (and, of course, R10 and R11) are preserved
 016A 682 : by this routine
 016A 683 :
 016A 684 : Assumptions:
 016A 685 :
 016A 686 : This routine makes two important assumptions.
 016A 687 :
 016A 688 : 1. If both of the input bytes contain only legal decimal digits, then
 016A 689 : it is only necessary to subtract 100 at most once to put all
 016A 690 : possible sums in the range 0..99. That is,
 016A 691 :
 016A 692 : $99 + 99 + 1 = 199 \text{ LSS } 200$
 016A 693 :
 016A 694 : 2. The result will be checked in some way to determine whether the
 016A 695 : result is nonzero so that the Z-bit can have its correct setting.
 016A 696 :-
 016A 697 :
 016A 698 ADD_PACKED_BYTE_STRING:
 016A 699 :
 56 71 9A 016A 700 MARK POINT ADD_SUB_BSBW_24
 56 71 9A 016A 701 MOVZBL -(R1),R6 ; Get byte from first string
 57 73 9A 016D 702 MARK POINT ADD_SUB_BSBW_24
 57 73 9A 016D 703 MOVZBL -(R3),R7 ; Get byte from second string
 0170 704 :
 0170 705 VAX\$ADD_PACKED_BYTE_R6_R7:: ; ASHP also uses this routine
 56 0000'CF46 90 0170 706 ADD_PACKED_BYTE_R6_R7:
 56 0000'CF46 90 0170 707 MOVB DECIMALSPACKED_TO_BINARY_TABLE[R6],-
 57 0000'CF47 90 0176 708 R6 ; Convert digits to binary
 57 0000'CF47 90 0176 709 MOVB DECIMALSPACKED_TO_BINARY_TABLE[R7],-
 57 56 80 017C 710 R7 ; Convert digits to binary
 57 58 80 017C 711 ADDB R6,R7 ; Form their sum
 57 58 80 017F 712 ADDB R8,R7 ; Add CARRY from last step
 63 8F 57 91 0182 713 CLRB R8 ; Assume no CARRY this time
 63 8F 57 91 0184 714 CMPB R7,#99 ; Check for CARRY
 07 1B 0188 715 BLEQU 10\$; Branch if within bounds
 58 01 90 018A 716 MOVB #1,R8 ; Propagate CARRY to next step
 57 64 8F 82 018D 717 SUBB #100,R7 ; Put R7 into interval 0..99
 75 0000'CF47 90 0191 718 10\$: MOVB DECIMALSBINARY_TO_PACKED_TABLE[R7],-
 0197 719 -(R5) ; Store converted sum byte
 05 0197 720 RSB

0198 722
0198 723 :+
0198 724 : Functional Description:
0198 725 :
0198 726 : This routine takes two packed decimal strings whose descriptors
0198 727 : are passed as input parameters, subtracts one string from the
0198 728 : other, and places their sum into another (perhaps identical)
0198 729 : packed decimal string.
0198 730 :
0198 731 : At the present time, the result is placed into a 16-byte storage
0198 732 : area while the difference is being evaluated. This drastically reduces
0198 733 : the number of different cases that must be dealt with as each
0198 734 : pair of bytes in the two input strings is added.
0198 735 :
0198 736 : The signs of the two input strings have already been dealt with so
0198 737 : this routine performs subtraction in all cases, even if the original
0198 738 : entry was at ADDP4 or ADDP6.
0198 739 :
0198 740 : Input Parameters:
0198 741 :
0198 742 : R0<4:0> - Number of digits in first input decimal string
0198 743 : R1 - Address of least significant digit of first input
0198 744 : decimal string (the byte containing the sign)
0198 745 :
0198 746 : R2<4:0> - Number of digits in second input decimal string
0198 747 : R3 - Address of least significant digit of second input
0198 748 : decimal string (the byte containing the sign)
0198 749 :
0198 750 : R4<4:0> - Number of digits in output decimal string
0198 751 : R5 - Address of one byte beyond least significant digit of
0198 752 : intermediate string stored on the stack
0198 753 :
0198 754 : R6<3:0> - Sign of first input string in preferred form
0198 755 : R7<3:0> - Sign of second input string in preferred form
0198 756 :
0198 757 : R11 - Saved PSL (Z-bit is set, other condition codes are clear)
0198 758 :
0198 759 : (SP) - Saved R5, address of least significant digit of ultimate
0198 760 : destination string.
0198 761 : 4(SP) - Beginning of 20-byte buffer to hold intermediate result
0198 762 :
0198 763 : Output Parameters:
0198 764 :
0198 765 : The particular input operation (ADDPx or SUBPx) is completed in
0198 766 : this routine. See the routine headers for the four routines that
0198 767 : request addition or subtraction for a list of output parameters
0198 768 : from this routine.
0198 769 :
0198 770 : Algorithm for Choice of Sign:
0198 771 :
0198 772 : The choice of sign for the output string is not nearly so
0198 773 : straightforward as it is in the case of addition. One approach that is
0198 774 : often taken is to make a reasonable guess at the sign of the result.
0198 775 : If the final subtraction causes a BORROW, then the choice was incorrect.
0198 776 : The sign must be changed and the result must be replaced by its tens
0198 777 : complement.
0198 778 :

.SUBTITLE SUBTRACT_PACKED - Subtract Two Packed Decimal Strings

Functional Description:

This routine takes two packed decimal strings whose descriptors are passed as input parameters, subtracts one string from the other, and places their sum into another (perhaps identical) packed decimal string.

At the present time, the result is placed into a 16-byte storage area while the difference is being evaluated. This drastically reduces the number of different cases that must be dealt with as each pair of bytes in the two input strings is added.

The signs of the two input strings have already been dealt with so this routine performs subtraction in all cases, even if the original entry was at ADDP4 or ADDP6.

Input Parameters:

R0<4:0> - Number of digits in first input decimal string
R1 - Address of least significant digit of first input decimal string (the byte containing the sign)
R2<4:0> - Number of digits in second input decimal string
R3 - Address of least significant digit of second input decimal string (the byte containing the sign)
R4<4:0> - Number of digits in output decimal string
R5 - Address of one byte beyond least significant digit of intermediate string stored on the stack
R6<3:0> - Sign of first input string in preferred form
R7<3:0> - Sign of second input string in preferred form
R11 - Saved PSL (Z-bit is set, other condition codes are clear)
(SP) - Saved R5, address of least significant digit of ultimate destination string.
4(SP) - Beginning of 20-byte buffer to hold intermediate result

Output Parameters:

The particular input operation (ADDPx or SUBPx) is completed in this routine. See the routine headers for the four routines that request addition or subtraction for a list of output parameters from this routine.

Algorithm for Choice of Sign:

The choice of sign for the output string is not nearly so straightforward as it is in the case of addition. One approach that is often taken is to make a reasonable guess at the sign of the result. If the final subtraction causes a BORROW, then the choice was incorrect. The sign must be changed and the result must be replaced by its tens complement.

0198 779 :
 0198 780 : This routine does not guess. Instead, it chooses the input string of
 0198 781 : the larger absolute magnitude as the minuend for this internal
 0198 782 : routine and chooses its sign as the sign of the result.
 0198 783 : This algorithm is actually more efficient than the reasonable
 0198 784 : guess method and is probably better than a guess method that is never
 0198 785 : wrong. All complete bytes that are processed in the sign evaluation
 0198 786 : preprocessing loop are eliminated from consideration in the
 0198 787 : subtraction loop, which has a higher cost per byte.
 0198 788 :
 0198 789 : The actual algorithm is as follows. (Note that both input strings have
 0198 790 : already had leading zeros stripped so their lengths reflect
 0198 791 : significant digits.)
 0198 792 :
 0198 793 : 1. If the two strings have unequal lengths, then choose the sign of
 0198 794 : the string that has the longer length.
 0198 795 :
 0198 796 : 2. For strings of equal length, choose the sign of the string whose
 0198 797 : most significant byte is larger in magnitude.
 0198 798 :
 0198 799 : 3. If the most significant bytes test equal, then decrease the
 0198 800 : lengths of each string by one byte, drop the previous most
 0198 801 : significant bytes, and go back to step 2.
 0198 802 :
 0198 803 : 4. If the two strings test equal, it is not necessary to do any
 0198 804 : subtraction. The result is identically zero.
 0198 805 :
 0198 806 : Note that the key to this routine's efficiency is that high order
 0198 807 : bytes that test equal in this loop are dropped from consideration in
 0198 808 : the more complicated subtraction loop.
 0198 809 :-

0198 810 SUBTRACT PACKED:
 50 04 01 EF 0198 811 EXTZV #1,#4,R0,R0 ; Convert digit count to byte count
 52 04 01 EF 019D 812 EXTZV #1,#4,R2,R2 ; Do it for both strings
 52 50 D1 01A2 813 CMPL R0,R2 ; We want to compare the byte counts
 3C 1F 01A5 814 BLSSU 40\$; R0/R1 represent the smaller string
 2A 1A 01A7 815 BGTRU 30\$; R2/R3 represent the smaller string
 01A9 816 :
 01A9 817 : The two input strings have an equal number of bytes. Compare magnitudes to
 01A9 818 : determine which string is really larger. If the two strings test equal, then
 01A9 819 : skip the entire subtraction loop.
 01A9 820 :
 58 51 50 C3 01A9 821 SUBL3 R0,R1,R8 ; Point R8 to low address end of R0/R1
 59 53 52 C3 01AD 822 SUBL3 R2,R3,R9 ; Point R9 to low address end of R2/R3
 50 D5 01B1 823 TSTL R0 ; See if both strings have zero bytes
 OC 13 01B3 824 BEQL 20\$; Still need to check low order digit
 01B5 825 :
 89 88 91 01B5 826 MARK_POINT ADD_SUB_24
 29 1F 01B8 827 10\$: CMPB (R8)+,(R9)+ ; Compare most significant bytes
 17 1A 01BA 828 BLSSU 40\$; R0/R1 represent the smaller string
 52 D7 01BC 829 BGTRU 30\$; R2/R3 represent the smaller string
 F4 50 F5 01BE 830 DECL R2 ; Keep R2 in step with R0
 01C1 831 SOBGTR R0,10\$; ... which gets decremented here
 01C1 832 : At this point, we have reduced both input strings to single bytes that

01C1 834 : contain a sign "digit" and may contain a digit in the high order nibble
 01C1 835 : if the original digit counts were nonzero.
 01C1 836
 01C1 837 MARK POINT ADD_SUB 24
 58 68 OF 8B 01C1 R78 20\$: BICB3 #^B00001111,(R8),R8 : Look only at digit, ignoring sign
 59 69 OF 8B 01C5 339 MARK POINT ADD SUB 24
 59 58 91 01C9 840 BICB3 #^B00001111,(R9),R9 : Get the digit from the other string
 15 1F 01CC 841 CMPB R8,R9 : Compare these digits
 03 1A 01CE 842 BLSSU 40\$: R0/R1 represent the smaller string
 01D0 843 BGTRU 30\$: R2/R3 represent the smaller string
 01D0 844
 01D0 845 : The two strings have identical magnitudes. Enter the end processing code
 01D0 846 : with the intermediate result unchanged (that is, zero).
 01D0 847
 FF6A 31 01D0 848 BRW ADD_SUBTRACT_EXIT ; Join the common completion code
 01D3 849
 01D3 850 : The string described by R0 and R1 has the larger magnitude. Choose its sign.
 01D3 851 : Then swap the two string descriptors so that the main subtraction loops
 01D3 852 : always have R2 and R3 describing the larger string. Note that the use of
 C1D3 853 : R6 and R7 as scratch leaves R7<31:8> in an UNPREDICTABLE state.
 01D3 854
 59 56 90 01D3 855 30\$: MOVB R6,R9 : Load preferred sign into R9
 56 50 7D 01D6 856 MOVQ R0,R6 : Save the longer
 50 52 7D 01D9 857 MOVQ R2,R0 : Store the shorter on R0 and R1
 52 56 7D 01DC 858 MOVQ R6,R2 : ... and store the longer in R2 and R3
 57 D4 01DF 859 CLRL R7 : Insure that R7<31:8> is zero
 03 11 01E1 860 BRB 50\$: Continue along common code path
 01E3 861
 01E3 862 : The string described by R2 and R3 has the larger magnitude. Choose its sign.
 01E3 863
 59 57 90 01E3 864 40\$: MOVB R7,R9 : Load preferred sign into R9
 01E6 865
 52 50 C2 01E6 866 50\$: SUBL R0,R2 : Make R2 a difference (R2 GEQU 0)
 03 59 E9 01E9 867 BLBC R9,60\$: Check if sign is negative
 58 08 88 01EC 868 BISB #P\$LSM_N,R11 : ... so the saved N-bit can be set
 01EF 869
 01EF 870 MARK POINT ADD_SUB 24
 56 61 OF 8B 01EF 871 60\$: BICB3 #^B00001111,(R1),R6 : Get least significant digit to R6
 01F3 872 MARK POINT ADD SUB 24
 57 63 OF 8B 01F3 873 BICB3 #^B00001111,(R3),R7 : Get least significant digit to R7
 58 D4 01F7 874 CLRL R8 : Start subtracting with BORROW off
 0032 30 01F9 875 BSBW SUB_PACKED_BYT_R6_R7 : Subtract the two low order digits
 01FC 876
 01FC 877 : R0 contains the number of bytes remaining in the smaller string
 01FC 878 : R2 contains the difference in bytes between the two input strings
 01FC 879
 50 D5 01FC 880 TSTL R0 : Does smaller string have any room?
 06 13 01FE 881 BEQL 80\$: Skip loop if no room at all
 0200 882
 0025 30 0200 883 70\$: BSBW SUB_PACKED_BYT_STRING : Subtract the next two bytes
 FA 50 F5 0203 884 SOBGTR R0,70\$: Check for end of loop
 0206 885
 52 D5 0206 886 80\$: TSTL R2 : Does one of the strings have more?
 16 13 0208 887 BEQL 110\$: Skip next loops if all done
 020A 888
 0D 58 E9 020A 889 90\$: BLBC R8,100\$: Life is simple if BORROW clear
 020D 890

56	D4	020D	891	CLRL R6		; Otherwise, BORROW must propagate
57	73	9A	020F	892	MARK POINT	ADD_SUB_24
0019		30	0212	893	MOVZBL -(R3), R7	; So subtract BORROW from single string
F2	52	F5	0215	894	BSBW SUB_PACKED_BYTE_R6_R7	; Use the special entry point
06		11	0218	895	S0BGTR R2, 90S	; Check for this string exhausted
			021A	896	BRB 110S	; Join common completion code
75	73	90	021A	897	100\$: MARK_POINT	ADD_SUB_24
FA	52	F5	021D	898	MOVB -(R3), -(R5)	; Simply move src to dst if no BORROW
			0220	899	S0BGTR R2, 100S	; ... until we're all done
			0220	900		
			0220	901		
			0220	902	110\$:	
			0220	903		
			0220	904	::::: ***** BEGIN TEMP *****	
			0220	905	::::: THE FOLLOWING HALT INSTRUCTION SHOULD BE REPLACED WITH THE CORRECT	
			0220	906	ABORT CODE.	
			0220	907	::::: THE HALT IS SIMILAR TO THE	
			0220	908	MICROCODE CANNOT GET HERE	
			0220	909	::::: ERRORS THAT OTHER IMPLEMENTATIONS USE.	
58	D5	0220	910	tstl r8		; If BORROW is set here, we blew it
01	13	0222	911	beql 120S		; Branch out if OK
00		0224	912	halt		; This will cause an OPCDEC exception
		0225	913	120\$:		
		0225	914		::::: ***** END TEMP *****	
FF15	31	0225	915	BRW ADD_SUBTRACT_EXIT		; Join common completion code
		0225	916			
		0225	917			
		0225	918			
		0225	919			
		0225	920			
		0225	921			

0228 924 .SUBTITLE SUB_PACKED_BYTE - Subtract Two Bytes Containing Decimal Digi
0228 925 :+
0228 926 : Functional Description:
0228 927 :
0228 928 : This routine takes as input two bytes containing decimal digits and
0228 929 : produces a byte containing their difference. This result is stored in
0228 930 : the output string. Each of the input bytes is converted to a binary
0228 931 : number (with a table-driven conversion), the first number is
0228 932 : subtracted from the second, and the difference is converted back to
0228 933 : two decimal digits stored in a byte.
0228 934 :
0228 935 : This routine makes no provisions for bytes that contain illegal
0228 936 : decimal digits. We are using the UNPREDICTABLE statement in the
0228 937 : architectural description of the decimal instructions to its fullest.
0228 938 :
0228 939 : The bytes that contain a pair of packed decimal digits can either
0228 940 : exist in packed decimal strings located by R1 and R3 or they can
0228 941 : be stored directly in registers. In the former case, the digits must
0228 942 : be extracted from registers before they can be used in later operations
0228 943 : because the difference will be used as an index register.
0228 944 :
0228 945 : For entry at SUB_PACKED_BYTE_STRING:
0228 946 :
0228 947 : Input Parameters:
0228 948 :
0228 949 : R1 - Address one byte beyond byte containing subtrahend
0228 950 : R3 - Address one byte beyond byte containing minuend
0228 951 : R5 - Address one byte beyond location to store difference
0228 952 :
0228 953 : R8 - BORROW from previous byte (R8 is either 0 or 1)
0228 954 :
0228 955 : Implicit Input:
0228 956 :
0228 957 : R6 - Scratch
0228 958 : R7 - Scratch
0228 959 :
0228 960 : Output Parameters:
0228 961 :
0228 962 : R1 - Decreased by one to point to current byte
0228 963 : in subtrahend string
0228 964 : R3 - Decreased by one to point to current byte
0228 965 : in minuend string
0228 966 : R5 - Decreased by one to point to current byte
0228 967 : in difference string
0228 968 :
0228 969 : R8 - Either 0 or 1, reflecting whether this most recent
0228 970 : subtraction resulted in a BORROW from the next byte.
0228 971 :
0228 972 : For entry at SUB_PACKED_BYTE_R6_R7:
0228 973 :
0228 974 : Input Parameters:
0228 975 :
0228 976 : R6<7:0> - Byte containing decimal digit pair for subtrahend
0228 977 : R6<31:8> - MBZ
0228 978 : R7<7:0> - Byte containing decimal digit pair for minuend
0228 979 : R7<31:8> - MBZ
0228 980 :

```

0228 981 : R5 - Address one byte beyond location to store difference
0228 982 : 
0228 983 : R8 - BORROW from subtraction of previous byte
0228 984 : (R8 is either 0 or 1)
0228 985 : 
0228 986 : Output Parameters:
0228 987 : 
0228 988 : R5 - Decreased by one to point to current byte
0228 989 : in difference string
0228 990 : 
0228 991 : R8 - Either 0 or 1, reflecting whether this most recent
0228 992 : subtraction resulted in a BORROW from the next byte.
0228 993 : 
0228 994 : Side Effects:
0228 995 : 
0228 996 : R6 and R7 are modified by this routine
0228 997 : 
0228 998 : R0, R2, R4, and R9 (and, of course, R10 and R11) are preserved
0228 999 : by this routine
0228 1000 : 
0228 1001 : Assumptions:
0228 1002 : 
0228 1003 : This routine makes two important assumptions.
0228 1004 : 
0228 1005 : 1. If both of the input bytes contain only legal decimal digits, then
0228 1006 : it is only necessary to add 100 at most once to put all
0228 1007 : possible differences in the range 0..99. That is,
0228 1008 : 
0228 1009 : 0 - 99 - 1 = -100
0228 1010 : 
0228 1011 : 2. The result will be checked in some way to determine whether the
0228 1012 : result is nonzero so that the Z-bit can have its correct setting.
0228 1013 :-
0228 1014 : 
0228 1015 SUB_PACKED_BYTE_STRING:
0228 1016 : 
0228 1017 MARK POINT ADD_SUB_BSBW_24
56 71 9A 0228 1018 MOVZBL -(R1),R6 ; Get byte from first string
0228 1019 MARK POINT ADD_SUB_BSBW_24
57 73 9A 0228 1020 MOVZBL -(R3),R7 ; Get byte from second string
0228 1021 : 
0228 1022 SUB_PACKED_BYTE_R6,R7:
56 0000'CF46 90 022E 1023 MOVB DEIMAL$PACKED_TO_BINARY_TABLE[R6],-
0234 1024 R6 ; Convert digits to binary
57 0000'CF47 90 0234 1025 MOVB DEIMAL$PACKED_TO_BINARY_TABLE[R7],-
023A 1026 R7 ; Convert digits to binary
57 56 82 023A 1027 SUBB R6,R7 ; Form their difference
57 58 82 023D 1028 SUBB R8,R7 ; Include BORROW from last step
04 19 0240 1029 BLSS 10$ ; Branch if need to BORROW
58 94 0242 1030 CLR B R8 ; No BORROW next time
07 11 0244 1031 BRB 20$ ; Join common exit code
0246 1032 : 
57 64 BF 80 0246 1033 10$: ADDB #100,R7 ; Put R7 into interval 0..99
58 01 90 024A 1034 MOVB #1,R8 ; Propogate BORROW to next step
024D 1035 : 
75 0000'CF47 90 024D 1036 20$: MOVB DEIMAL$BINARY_TO_PACKED_TABLE[R7],-
0253 1037 -(RS) ; Store converted sum byte

```

VAX\$DECIMAL_ARITHMETIC
V04-001

N 8
- VAX-11 Packed Decimal Arithmetic Instr 8-JAN-1985 17:27:01 VAX/VMS Macro V04-00
SUB_PACKED_BYT^E - Subtract Two Bytes Con 5-SEP-1984 00:44:34 [EMULAT.BUGSRC]VAXARITH.MAR;1 Page 23
(11)
05 0253 1038 RSB

0254 1040 .SUBTITLE STORE_RESULT - Store Decimal String
 0254 1041 ::
 0254 1042 :: Functional Description:
 0254 1043 ::
 0254 1044 :: This routine takes a packed decimal string that typically contains
 0254 1045 :: the result of an arithmetic operation and stores it in another
 0254 1046 :: decimal string whose descriptor is specified as an input parameter
 0254 1047 :: to the original arithmetic operation.
 0254 1048 ::
 0254 1049 :: The string is stored from the high address end (least significant
 0254 1050 :: digits) to the low address end (most significant digits). This order
 0254 1051 :: allows all of the special cases to be handled in the simplest fashion.
 0254 1052 ::
 0254 1053 :: Input Parameters:
 0254 1054 ::
 0254 1055 :: R1 - Address one byte beyond high address end of input string
 0254 1056 :: (Note that this string must be at least 17 bytes long.)
 0254 1057 ::
 0254 1058 :: R4<4:0> - Number of digits in ultimate destination
 0254 1059 :: R5 - Address one byte beyond destination string
 0254 1060 ::
 0254 1061 :: R11 - Contains saved condition codes
 0254 1062 ::
 0254 1063 ::
 0254 1064 :: Implicit Input:
 0254 1065 :: The input string must be at least 17 bytes long to contain a potential
 0254 1066 :: carry out of the highest digit when doing an add of two large numbers.
 0254 1067 :: This carry out of the last byte will be detected and reported as a
 0254 1068 :: decimal overflow, either as an exception or simply by setting the V-bit.
 0254 1069 ::
 0254 1070 :: The least significant digit (highest addressed byte) cannot contain a
 0254 1071 :: sign digit because that would cause the Z-bit to be incorrectly cleared.
 0254 1072 ::
 0254 1073 :: Output Parameters:
 0254 1074 ::
 0254 1075 :: R11<PSL\$V_Z> - Cleared if a nonzero digit is stored in output string
 0254 1076 :: R11<PSL\$V_V> - Set if a nonzero digit is detected after the output
 0254 1077 :: string is exhausted
 0254 1078 ::
 0254 1079 :: A portion of the result (dictated by the size of R4 on input) is
 0254 1080 :: moved to the destination string.
 0254 1081 ::-
 0254 1082 ::
 0254 1083 :: STORE_RESULT:
 54 FF 54 D6 0254 1084 INCL R4 ; Want number of "complete" bytes in
 8F 78 0256 1085 ASHL #1,R4,R0 ; output string
 50 025A 1086 BEQL 30\$; Skip first loop if none
 0B 13 025B 1087
 025D 1088 MARK_POINT ADD_SUB_BSBW_24
 75 71 90 025D 1089 10\$: MOVBL -(R1),-(R5) ; Move the next complete byte
 03 13 0260 1090 BEQL 20\$; Check whether to clear Z-bit
 5B 04 8A 0262 1091 BICB #PSLSM_Z,R11 ; Clear Z-bit if nonzero
 F5 50 F5 0265 1092 20\$: SOBGTR R0,10\$; Keep going?
 10 54 E9 0268 1093
 026B 1094 30\$: BLBC R4,50\$; Was original R4 odd? Branch if yes
 MARK_POINT ADD_SUB_BSBW_24

71 F0 8F 8B 026B 1096 BICB3 #^B11110000,-(R1),-(R5) ; If R4 was even, store half a byte
 75 026F
 03 13 0270 1097 BEQL 40\$
 5B 04 8A 0272 1098 BICB #PSLSM_Z,R11 : Need to check for zero here, too
 0275 1099 MARK_POINT ADD_SUB_BSBW_24 : Clear Z-bit if nonzero
 61 F0 8F 93 0275 1100 40\$: BITB #^B11110000,(R1) : If high order nibble is nonzero,
 13 12 0279 1101 BNEQ 70\$; ... then overflow has occurred
 027B 1102
 027B 1103 : The entire destination has been stored. We must now check whether any of
 027B 1104 : the remaining input string is nonzero and set the V-bit if nonzero is
 027B 1105 : detected. Note that at least one byte of the output string has been examined
 027B 1106 : in all cases already. This makes the next byte count calculation correct.
 027B 1107
 54 04 54 D7 027B 1108 50\$: DECL R4 ; Restore R4 to its original self
 01 EF 027D 1109 EXTZV #1,#4,R4,R0 ; Extract a byte count
 50 0281
 50 10 50 83 0282 1110 SUBB3 R0,#16,R0 ; Loop count is 16 minus byte count
 0286 1111
 0286 1112 : Note that the loop count can never be zero because we are testing a 17-byte
 0286 1113 : string and the largest output string can be 16 bytes long.
 0286 1114
 71 95 0286 1115 MARK_POINT ADD_SUB_BSBW_24
 04 12 0288 1116 60\$: TSTB -(R1) ; Check next byte for nonzero
 F9 50 F5 028A 1117 BNEQ 70\$; Nonzero means overflow has occurred
 028D 1118 SOBGTR R0,60\$; Check for end of this loop
 05 028D 1120 RSB ; This is return path for no overflow
 028E 1121
 5B 02 88 028E 1122 70\$: BISB #PSLSM_V,R11 ; Indicate that overflow has occurred
 05 0291 1123 RSB ; ... and return to the caller

:LJK0045	0292	1
:LJK0045	0292	2
:LJK0045	0292	3
:LJK0045	0292	4
:LJK0045	0292	5
:LJK0045	0292	6
:LJK0045	0292	7
:LJK0045	0292	8
:LJK0045	0292	9
:LJK0045	0292	10
:LJK0045	0292	11
:LJK0045	0292	12
:LJK0045	0292	13
:LJK0045	0292	14
:LJK0045	0292	15
:LJK0045	0292	16
:LJK0045	0292	17
:LJK0045	0292	18
:LJK0045	0292	19
:LJK0045	0292	20
:LJK0045	0292	21
:LJK0045	0292	22
:LJK0045	0292	23
:LJK0045	0292	24
:LJK0045	0292	25
:LJK0045	0292	26
:LJK0045	0292	27
:LJK0045	0292	28
:LJK0045	0292	29
:LJK0045	0292	30
:LJK0045	0292	31
:LJK0045	0292	32
:LJK0045	0292	33
:LJK0045	0292	34
:LJK0045	0292	35
:LJK0045	0292	36
:LJK0045	0292	37
:LJK0045	0292	38
:LJK0045	0292	39
:LJK0045	0292	40
:LJK0045	0292	41
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:LJK0045	0292	44
:LJK0045	0292	45
:LJK0045	0292	46
:LJK0045	0292	47
:LJK0045	0292	48
:LJK0045	0292	49
:LJK0045	0292	50
:LJK0045	0292	51
:LJK0045	0292	52
:LJK0045	0292	53
:LJK0045	0292	54
:LJK0045	0292	55
:LJK0045	0292	56
:LJK0045	0292	57

.SUBTITLE - CHECK_WRITE_ACCESS - Check Writability of Decimal String

Functional Description:

The ADDP4 and SUBP4 instructions are unique in that they are the only two instructions that read and write the same packed decimal string. They are, in fact, implemented as ADDP6 and SUBP6 where the second input string, ADDend2 or MINuend, and the result string, SUM or DIFFERENCE, are the same. But ADDP6 and SUBP6, as well as all other packed decimal instructions except ADDP4 and SUBP4, produce UNPREDICTABLE results when an output string overlaps any input string. With this interpretation, ADDP4 and SUBP4 are the only packed decimal instructions that permit overlapping packed decimal strings. (Note that the result string, SUM or DIFFERENCE, may not overlap the first input string, ADDend or SUBtrahend.)

The implementation of ADDP4 and SUBP4, interpreted as ADDP6 and SUBP6 with overlapping strings, needs to protect itself from modifying memory until the entire instruction can execute to completion. Otherwise, the instruction will be restarted with a different initial state than it first had. This is accomplished by probing the result string for write access before execution begins. This routine receives control if that PROBEW fails.

In the interest of simplicity, the PROBEW that is executed always uses 16 as the byte count of the result string. This routine can then be called not only when the output string is inaccessible but also when

The output string is shorter than 16 bytes.

The address of the output string is within 16 bytes of the end of the page.

The page containing the string is writable.

The next page is not writable.

This routine distinguishes inaccessible strings from this rare case of a PROBEW failure.

In other words, this routine checks the write accessibility of packed decimal strings. If the entire string is writable, control is passed back to the caller where emulation continues. If the string is not writable, an access violation is generated.

Input Parameters:

R4 - Digit count of result string

R5 - Address of most significant digit in result string

Output Parameters:

None

Implicit Output:

If the string with its correct byte count is writable, control

:LJK0045 0292 .58 : is returned to the caller.

:LJK0045 0292 .59 : If the string with its correct digit count is not writable, then the

:LJK0045 0292 .60 : string is accessed to force an access violation to occur. This will

:LJK0045 0292 .61 : cause control to be transferred eventually to the access violation

:LJK0045 0292 .62 : handler at the end of this module that will back out the instruction

:LJK0045 0292 .63 : so that it can be restarted.

:LJK0045 0292 .64 : Assumption:

:LJK0045 0292 .65 : This routine assumes that the check for illegal digit count has

:LJK0045 0292 .66 : already been made so that R4 is between 0 and 31 inclusive.

:LJK0045 0292 .67 : Note:

:LJK0045 0292 .68 : It is necessary to actually touch the inaccessible page so that the

:LJK0045 0292 .69 : correct reason mask is generated. Logic that uses the PROBEW

:LJK0045 0292 .70 : instruction can determine the inaccessible virtual address but is

:LJK0045 0292 .71 : unable to distinguish, for example, between length violations and page

:LJK0045 0292 .72 : protection violations.

:LJK0045 0292 .73 :--

:LJK0045 0292 .74 :.80 CHECK_WRITE_ACCESS:

54 04 54 DD 0292 .81 PUSHL R4 ; Save the digit count

:LJK0045 54 01 EF 0294 .82 EXTZV #1,#4,R4,R4 ; Convert digits to bytes

:LJK0045 65 54 54 D6 0299 .83 INCL R4 ; Count the byte that contains the sign

:LJK0045 65 54 00 0D 029B .84 PROBEW #0,R4,(R5) ; Check writability of the string

:LJK0045 04 13 029F .85 BEQL 10\$; Branch if string cannot be written

:LJK0045 54 8E D0 02A1 .86 MOVL (SP)+,R4 ; Restore saved R4

:LJK0045 05 02A4 .87 RSB ; Return to caller. String is OK

:LJK0045 02A5 .88 : The first and last bytes of the string are touched (written) with an

:LJK0045 02A5 .89 ADDB2 #0,xxx

:LJK0045 02A5 .90 : instruction. This instruction causes the correct access violation reason

:LJK0045 02A5 .91 mask to be generated but does not modify the contents of locations that are

:LJK0045 02A5 .92 accessible. Note that at least one of the following two ADDB2 instructions

:LJK0045 02A5 .93 is guaranteed to generate an access violation, transferring control in a

:LJK0045 02A5 .94 rather complicated way to the ADD_SUB_BSBW_4 access violation handler.

:LJK0045 02A5 .95 : MARK_POINT ADD_SUB_BSBW_4

:LJK0045 65 00 80 02A5 .96 ADDB2 #0,(R5) ; Touch the first byte

:LJK0045 02A8 .97 MARK_POINT ADD_SUB_BSBW_4

:LJK0045 6544 00 80 02A8 .98 ADDB2 #0,(R5)[R4] ; Touch the last byte

:LJK0045 02AC .99 : We should never reach here unless the PROBE instruction is broken. We

:LJK0045 02AC .100 10\$: will leave this code path in for now but remove it at the same time we

:LJK0045 02AC .101 change the other two HALT instructions in this module into software

:LJK0045 02AC .102 generated machine checks exceptions.

:LJK0045 02AC .103 :.104 ***** BEGIN TEMP *****

:LJK0045 02AC .105 : THE FOLLOWING HALT INSTRUCTION SHOULD BE REPLACED WITH THE CORRECT

:LJK0045 02AC .106 ABORT CODE.

:LJK0045 02AC .107 :.108

:LJK0045 02AC .109 :.110

:LJK0045 02AC .111 :.112

:LJK0045 02AC .113 :.114

VAX\$DECIMAL_ARITHMETIC
V04-001

F 9
- VAX-11 Packed Decimal Arithmetic Instr 8-JAN-1985 17:27:01 VAX/VMS Macro V04-00 Page 28
- CHECK_WRITE_ACCESS - Check Writability 5-SEP-1984 00:44:34 [EMULAT.BUGSRC]VAXARITH.MAR;1 (13)
;LJK0045
;LJK0045
;LJK0045
00 02AC .114 halt ; This will cause an OPCDEC exception
02AD .115 ;;
02AD .116 ;;; ***** END TEMP *****

02AD 1125
 02AD 1126
 02AD 1127
 02AD 1128
 02AD 1129
 02AD 1130
 02AD 1131
 02AD 1132
 02AD 1133
 02AD 1134
 02AD 1135
 02AD 1136
 02AD 1137
 02AD 1138
 02AD 1139
 02AD 1140
 02AD 1141
 02AD 1142
 02AD 1143
 02AD 1144
 02AD 1145
 02AD 1146
 02AD 1147
 02AD 1148
 02AD 1149
 02AD 1150
 02AD 1151
 02AD 1152
 02AD 1153
 02AD 1154
 02AD 1155
 02AD 1156
 02AD 1157
 02AD 1158
 02AD 1159
 02AD 1160
 02AD 1161
 02AD 1162
 02AD 1163
 02AD 1164
 02AD 1165
 02AD 1166
 02AD 1167
 02AD 1168
 02AD 1169
 02AD 1170
 02AD 1171
 02AD 1172
 02AD 1173
 02AD 1174
 02AD 1175
 02AD 1176
 02AD 1177
 02AD 1178
 02AD 1179
 02AD 1180
 02AD 1181

.SUBTITLE VAXSMULP - Multiply Packed

Functional Description:

The multiplicand string specified by the multiplicand length and multiplicand address operands is multiplied by the multiplier string specified by the multiplier length and multiplier address operands. The product string specified by the product length and product address operands is replaced by the result.

Input Parameters:

R0 - mulrlen.rw	Number of digits in multiplier string
R1 - mulraddr.ab	Address of multiplier string
R2 - mulrlen.rw	Number of digits in multiplicand string
R3 - muldaddr.ab	Address of multiplicand string
R4 - prodlen.rw	Number of digits in product string
R5 - prodaddr.ab	Address of product string

Output Parameters:

R0 = 0
 R1 = Address of the byte containing the most significant digit of the multiplier string
 R2 = 0
 R3 = Address of the byte containing the most significant digit of the multiplicand string
 R4 = 0
 R5 = Address of the byte containing the most significant digit of the string containing the product

Condition Codes:

N <- product string LSS 0
 Z <- product string EQL 0
 V <- decimal overflow
 C <- 0

Register Usage:

This routine uses all of the general registers. The condition codes are computed at the end of the instruction as the final result is stored in the product string. R11 is used to record the condition codes.

Notes:

1. This routine uses a large amount of stack space to allow storage of intermediate results in a convenient form. Specifically, each digit pair of the longer input string is stored in binary in a longword on the stack. In addition, 32 longwords are set aside to hold the product intermediate result. Each longword contains a binary number between 0 and 99.

After the multiplication is complete, each longword is removed from the stack, converted to a packed decimal pair, and stored in the output string. Any nonzero cells remaining on the stack after the

	02AD	1182	:	output string has been completely filled are the indication of decimal overflow.
	02AD	1183	:	The purpose of this method of storage is to avoid decimal/binary or even byte/longword conversions during the calculation of intermediate results.
	02AD	1184	:	
	02AD	1185	:	
	02AD	1186	:	
	02AD	1187	:	
	02AD	1188	:	
	02AD	1189	:	
	02AD	1190	:	
	02AD	1191	:	
	02AD	1192	:	
	02AD	1193	:	
	02AD	1194	:	
	02AD	1195	VAX\$MULP::	
0FFF 8F BB	02AD	1196	PUSHR #^M<R0,R1,R2,R3,R4,R5,R6,R7,R8,R9,R10,R11>	; Save the lot
	02B1	1197		
	02B1	1198	ESTABLISH_HANDLER	- ; Store address of access
	02B1	1199	ARITH_ACCVIO	- ; violation handler
	02B6	1200		
	02B6	1201	ROPRAND_CHECK R4	; Insure that R4 is LEQU 31
	02C1	1202		
	02C1	1203	ROPRAND_CHECK R2	; Insure that R2 is LEQU 31
FD34' 30	02C9	1204	MARK_POINT MULP BSBW 0	
	02CC	1205	BSBW DECIMAL\$STRIP_ZEROS_R2_R3	; Strip high order zeros from R2/R3
	02CC	1206		
	02C9	1207	ROPRAND_CHECK R0	; Insure that R0 is LEQU 31
	02D4	1208	MARK_POINT MULP BSBW 0	
FD29' 30	02D4	1209	BSBW DECIMAL\$STRIP_ZEROS_R0_R1	; Strip high order zeros from R0/R1
	02D7	1210		
50 04 01 EF	02D7	1211	EXTZV #1,#4,R0,R0	; Convert digit count to byte count
50	02DB			
50	D6	02DC	1212 INCL R0	; Include least significant digit
52 04 01 EF	02DE	1213		
52	02E2			
52	D6	02E3	1214 EXTZV #1,#4,R2,R2	; Convert digit count to byte count
	02E5	1215 INCL R2	; Include least significant digit	
52	02E5	1216 CMPL R0,R2	; See which string is larger	
08	1A	02E8	1218 BGTRU 3S	; R2/R3 describes the longer string
58	52	02EA	1219 MOVQ R2,R8	; R8 and R9 describe the longer string
7E	50	02ED	1220 MOVQ R0,-(SP)	; Shorter string descriptor also saved
06	11	02F0	1221 BRB 6S	
58	50	7D	02F2 1222 3S: MOVQ R0,R8	; R8 and R9 describe the longer string
7E	52	7D	02F5 1224 MOVQ R2,-(SP)	; Shorter string descriptor also saved
	02F8	1225		
	02F8	1226 ; Create space for the output array on the stack (32 longwords of zeros)		
50 08 D0	02F8	1228 6S: MOVL #8,R0	; Eight pairs of quadwords	
	02FB	1229		
7E 7C	02FB	1230 10S: CLRQ -(SP)	; Clear one pair	
7E 7C	02FD	1231 CLRQ -(SP)	; ... and another	
F9 50 F5	02FF	1232 SOBGTR R0,10S	; Do all eight pairs	
57 5E D0	0302	1234 MOVL SP,R7	; Store beginning of output array in R7	
	0305	1235		
	0305	1236 ; The longer input array will be stored on the stack as an array of		

0305 1237 : longwords. Each array element contains a number between 0 and 99.
 0305 1238 : representing a pair of digits in the original packed decimal string.
 0305 1239 : Because the units digit is stored with the sign in packed decimal format,
 0305 1240 : it is necessary to shift the number as we store it. This is accomplished by
 0305 1241 : multiplying the number by ten.
 0305 1242 :
 0305 1243 : The longer array is described by R8 (byte count) and R9 (address of most
 0305 1244 : significant digit pair).
 0305 1245 :
 55 58 59 C1 0305 1246 ADDL3 R9,R8,R5 : Point R5 beyond sign digit
 54 58 D0 0309 1247 MOVL R8,R4 : R4 contains the loop count
 030C 1248 :
 030C 1249 : An array of longwords is allocated on the stack. R3 starts out pointing
 030C 1250 : at the longword beyond the top of the stack. The first remainder, guaranteed
 030C 1251 : to be zero, is "stored" here. The rest of the digit pairs are stored safely
 030C 1252 : below the top of the stack.
 030C 1253 :
 53 5E 53 58 CE 030C 1254 MNEGL R8,R3 : Stack grows toward lower addresses
 6E43 DE 030F 1255 MOVAL (SP)[R3],SP : Allocate the space
 53 5E 04 C3 0313 1256 SUBL3 #4,SP,R3 : Point R3 at next lower longword
 0317 1257 :
 51 0000'CF41 9A 0317 1258 MARK POINT MULP_R8
 1259 20\$: MOVZBL -(R5),R1 : Get next digit pair
 031A 1260 MOVZBL DECIMALSPACKED_TO_BINARY_TABLE[R1],-
 0320 1261 R1 : Convert digits to binary
 52 51 0A 7A 0320 1262 EMUL #10,R1,R2,R0 : Multiply by 10
 00000064 8F 7B 0325 1263 EDIV #100,R0,R2,(R3)+ : Divide by 100
 83 52 50 032B E6 54 F5 032E 1264 SOBGTR R4,20\$
 0331 1265 :
 63 52 D0 0331 1266 MOVL R2,(R3) : Store final quotient
 59 5E D0 0334 1267 MOVL SP,R9 : Remember array address in R9
 6E48 DF 0337 1268 PUSHAL (SP)[R8] : Store start of fixed size area
 033A 1269 :
 033A 1270 : Check for trailing zeros in the input array stored on the stack. If any are
 033A 1271 : present, they are removed and the product array is adjusted accordingly.
 033A 1272 :
 89 D5 033A 1273 30\$: TSTL (R9)+ : Is next number zero?
 08 12 033C 1274 BNEQ 40\$: Leave loop if nonzero
 57 04 C0 033E 1275 ADDL #4,R7 : Advance output pointer to next element
 F6 58 F5 0341 1276 SOBGTR R8,30\$: Keep going
 0344 1277 :
 0344 1278 : If we drop through the loop, then the entire input array is zero. There is
 0344 1279 : no need to perform any arithmetic because the product will be zero (and the
 0344 1280 : output array on the stack starts out as zero). The only remaining work is
 0344 1281 : to store the result in the output string and set the condition codes.
 0344 1282 :
 20 11 0344 1283 BRB 70\$: Exit to end processing
 0346 1284 :
 0346 1285 : Now multiply the input array by each successive digit pair. In order to
 0346 1286 : allow R10 to continue to locate ARITH ACCVIO while we execute this loop, it
 0346 1287 : is necessary to perform a small amount of register juggling. In essence,
 0346 1288 : R8 and R9 switch the identity of the string that they describe.
 0346 1289 :
 59 04 C2 0346 1290 40\$: SUBL #4,R9 : Readjust input array pointer
 7E 58 7D 0349 1291 MOVQ R8,-(SP) : Save R8/R9 descriptor on stack

58	08 AE	D0 034C	1292	MOVL 8(SP),R8	: Point R8 at start of 32-longword array
58	0080 C8	7D 0350	1293	MOVQ <32*4>(R8),R8	: Get descriptor that follows that array
59	58	C0 0355	1294	ADDL2 R8,R9	: Point R9 beyond sign byte
53	87	DE 0358	1296	50\$: MOVAL (R7)+,R3	: Output array address to R3
56	51 79	9A 035B	1297	MARK POINT MULP_AT_SP	
	0000'CF41	9A 035E	1298	MOVZBL -(R9),R1	: Next digit pair to R1
		0364	1300	MOVZBL DECIMALPACKED_TO_BINARY_TABLE[R1],-	
54	06 13	0364	1301	R6 BEQL 60\$: Convert digits to binary
	6E 7D	0366	1302	MOVA (SP),R4	: Skip the work if zero
	0104 30	0369	1303	BSBW EXTEND_STRING_MULTIPLY	: Input array descriptor to R4/R5
	E9 58	F5 036C	1304	60\$: SOBGTR R8,50\$: Do the work
		036F	1305		: Any more multiplier digits?
5E	08 CO	036F	1306	ADDL #8,SP	: Discard saved long string descriptor
5E	6E D0	0372	1307	70\$: MOVL (SP),SP	: Remove input array from stack
		0375	1309		
		0375	1310	: At this point, the product string is located in a 32-longword array on	
		0375	1311	: the top of the stack. Each longword corresponds to a pair of digits in	
		0375	1312	: the output string. As digits are removed from the stack, they are checked	
		0375	1313	: for nonzero to obtain the correct setting of the Z-bit. After the output	
		0375	1314	: string has been filled, the remainder of the product string is removed from	
		0375	1315	: the stack. If a nonzero result is detected at this stage, the V-bit is set.	
54	59 20	D0 0375	1317	MOVL #32,R9	: Set up array counter
	0098 CE	7D 0378	1318	MOVQ <<32*4> + ->>	: Skip over 32-longword array
		037D	1319	<2*4> + -	: and saved string descriptor
		037D	1320	<4*4> >(SP),R4	: to retrieve original R4 and R5

					.SUBTITLE Common Exit Path for VAX\$MULP and VAX\$DIVP
					:+ The code for VAX\$MULP and VAX\$DIVP merges at this point. The result is stored in an array of longwords at the top of the stack. The size of this array is stored in R9. The original R4 and R5 have been retrieved from the stack.
					Input Parameters:
					R4 - Contains byte count of destination string in R4 <1:4> R5 - Address of most significant digit of destination string R9 - Count of longwords in result array on stack
					Contents of result array
					Implicit Input:
					SIGNS OF TWO INPUT FACTORS (multiplier and multiplicand or divisor and dividend)
					MULTIPLY DIVIDE_EXIT:
04 00	5B 04	DC F0	037D 1343	MOVPSL R11	; Get current PSL
	5B		037F 1344	INSV #PSLSM_Z,#0,#4,R11	; Clear all codes except Z-bit
			0383		
			0384 1345	ESTABLISH_HANDLER -	; Store address of access
54 04	01 EF	0384 1346		ARITH ACCVIO	; violation handler again
	53	038D	0389 1347	EXTZV #1,#4,R4,R3	; Excess byte count to R3
57 55	3B 53	C1 13	038E 1348	BEQL 125\$; Skip to single digit code
55 57	01 C1	0390 1349		ADDL3 R3,R5,R7	; Remember address of sign byte
		0394 1350		ADDL3 #1,R7,R5	; Point R5 beyond end of product string
		0398 1351			
51 8E	D0 03	0398 1352	80\$: MOVL (SP)+,R1		; Remove next value from stack
	13	0398 1353	BEQL 90\$; Do not clear Z-bit if zero
5B 04	8A	039D 1354	BICB2 #PSLSM_Z,R11		; Clear Z-bit
		03A0 1355			
		03A0 1356	MARK_POINT MULP DIVP R9		
75 0000'CF41	90	03A0 1357	90\$: MOVB DECIMAL\$BINARY_TO_PACKED_TABLE[R1],-(R5)		
		03A6 1358			; Store converted sum byte
59	D7 03A6	1359	DECL R9		; One less element on the stack
1C	15 03A8	1360	BLEQ 116\$; Exit loop if result array exhausted
EB 53	F5 03AA	1361	S0BGTR R3,80\$; Keep going?
22 54	E9 03AD	1362			
		1363 100\$: BLBC R4,120\$; Different for even digit count
		03B0 1364			
		03B0 1365	: The output string consists of an odd number of digits. A complete digit		
		03B0 1366	: pair can be stored in the most significant (lowest addressed) byte of		
		03B0 1367	: the product string.		
		03B0 1368			
51 8E	D0 03B0	1369	MOVL (SP)+,R1		; Remove next value from stack
	13 03B3	1370	BEQL 110\$; Do not clear Z-bit if zero
5B 04	8A 03B5	1371	BICB2 #PSLSM_Z,R11		; Clear Z-bit
		03B8 1372			
		03B8 1373	MARK_POINT MULP DIVP R9		
75 0000'CF41	90	03B8 1374	110\$: MOVB DECIMAL\$BINARY_TO_PACKED_TABLE[R1],-(R5)		
		03BE 1375			; Store converted sum byte
59	D7 03BE	1376	DECL R9		; One less element on the stack

04 15 03C0 1377 BLEQ 116S ; Exit loop if result array exhausted
 38 11 03C2 1378 BRB 140S ; Perform overflow check
 03C4 1379
 03C4 1380 ; This loop executes if the result array has fewer elements than the output
 03C4 1381 ; string. The remaining bytes in the output string are filled with zeros.
 03C4 1382 ; There is no need for an overflow check.
 03C4 1383
 03C4 1384 MARK_POINT MULP_DIVP_8
 FB 75 94 03C4 1385 114S: CLRB -(R5) ; Store another zero byte
 53 F4 03C6 1386 116S: SOBGEQ R3,114S ; Any more room in output string
 03C9 1387
 38 11 03C9 1388 BRB 150S ; Determine sign of result
 03CB 1389
 03CB 1390 ; This code path is used in the case where the output digit count is 0 or 1.
 03CB 1391 ; R5 must be advanced
 03CB 1392
 57 55 D0 03CB 1393 125S: MOVL R5,R7 ; Remember address of output sign byte
 55 D6 03CE 1394 INCL R5 ; Advance R5 so common code can be used
 DB 11 03D0 1395 BRB 100S ; Join common code path
 03D2 1396
 03D2 1397 ; The output string consists of an even number of digits. Only the low order
 03D2 1398 ; nibble is stored in the most significant (lowest addresses) byte. A zero is
 03D2 1399 ; stored in the high order nibble. If the high order digit would have been
 03D2 1400 ; nonzero, the V-bit is set and the overflow check is bypassed because there
 03D2 1401 ; are faster ways to clean the stack if we do not have to check for nonzero
 03D2 1402 ; at the same time.
 03D2 1403
 51 51 BE D0 03D2 1404 120S: MOVL (SP)+,R1 ; Remove next value from stack
 51 0000'CF41 90 03D5 1405 MOVB DECIMALSBINARY_TO_PACKED_TABLE[R1],- ; Obtain converted sum byte
 03DB 1406
 03DB 1407 MARK_POINT MULP_DIVP_R9
 51 F0 8F 88 03DB 1408 BICB3 #^XF0,R1,-(R5) ; Store byte, clearing high order nibble
 75 03DF
 03 13 03E0 1409 BEQL 130S ; Do not clear Z-bit if zero
 51 5B 04 8A 03E2 1410 BICB2 #PSLSM_Z,R11 ; Clear Z-bit
 06 12 03E5 1411 130S: BITB #^XF0,R1 ; Is high order nibble nonzero?
 59 D7 03E9 1412 BNEQ 133S ; Yes, go set overflow bit
 D7 15 03EB 1413 DECL R9 ; One less element on the stack
 0B 11 03EF 1414 BLEQ 116S ; Exit loop if result array exhausted
 03F1 1415 BRB 140S ; Check rest of result array for nonzero
 03F1 1416
 03F1 1417 ; If we detect overflow, we need to adjust R9 to reflect the nonzero longword
 03F1 1418 ; removed from the stack before we enter the next code block that sets the
 03F1 1419 ; V-bit and cleans off the stack based on the contents of R9.
 03F1 1420
 59 D7 03F1 1421 133S: DECL R9 ; One more longword removed from stack
 03F3 1422
 03F3 1423 ; A nonzero digit has been discovered in a position that cannot be stored in
 03F3 1424 ; the output string. Set the V-bit, remove the rest of the product array from
 03F3 1425 ; the stack, and join the exit processing in the code that determines the sign
 03F3 1426 ; of the product.
 03F3 1427
 5B 02 88 03F3 1428 135S: BISB #PSLSM_V,R11 ; Set the overflow bit
 SE 6E49 DE 03F6 1429 MOVAL (SP)[R9],SP ; Clean off remaining product string
 07 11 03FA 1430 BRB 150S ; Go to code that determines the sign
 03FC 1431
 03FC 1432 ; The remainder of the product array must be removed from the stack. A nonzero

03FC 1433 : result causes the V-bit to be set and the rest of the loop to be skipped.
 03FC 1434 : Note that there is always a nonzero loop count remaining at this point.
 03FC 1435
 BE D5 03FC 1436 140\$: TSTL (SP)+
 F1 12 03FE 1437 BNEQ 133\$
 F9 59 F5 0400 1438 SOBGTR R9,140\$
 0403 1439
 0403 1440 : The final product string has been stored and the V- and Z-bits have their
 0403 1441 : correct settings. The sign of the product must be determined from the
 0403 1442 : signs of the two input strings. Opposite signs produce a negative product.
 0403 1443 : Same signs (in any representation) produce a plus sign in the output string.
 0403 1444
 SE 08 C0 0403 1445 150\$: ADDL #8,SP
 56 0C D0 0406 1446 MOVL #12,R6
 50 6E 7D 0409 1447 MOVQ (SP),R0
 50 04 01 EF 040C 1448 EXTZV #1,#4,R0,R0
 50 0410
 51 50 C0 0411 1449 ADDL R0,R1
 0414 1450 MARK POINT MULP DIVP 0
 61 F0 8F 8B 0414 1451 BICB3 #^B11110000,7R1),R0
 50 0418
 0419 1452
 0419 1453 CASE R0,TYPE=B,LIMIT=#10,<-
 0419 1454 220\$,-
 0419 1455 210\$,-
 0419 1456 220\$,-
 0419 1457 210\$,-
 0419 1458 220\$,-
 0419 1459 220\$,-
 0419 1460 220\$,-
 >
 0429 1461
 54 01 D0 0429 1462 210\$: MOVL #1,R4
 02 11 042C 1463 BRB 230\$
 042E 1464
 54 D4 042E 1465 220\$: CLRL R4
 0430 1466
 52 08 AE 7D 0430 1467 230\$: MOVQ 8(SP),R2
 52 04 01 EF 0434 1468 EXTZV #1,#4,R2,R2
 52 0438
 53 52 C0 0439 1469 ADDL R2,R3
 043C 1470 MARK POINT MULP DIVP 0
 63 F0 8F 8B 043C 1471 BICB3 #^B11110000,7R3),R2
 52 0440
 0441 1472
 0441 1473 CASE R2,TYPE=B,LIMIT=#10,<-
 0441 1474 250\$,-
 0441 1475 240\$,-
 0441 1476 250\$,-
 0441 1477 240\$,-
 0441 1478 250\$,-
 0441 1479 250\$,-
 0441 1480 250\$,-
 >
 0451 1481
 10 09 54 D6 0451 1482 240\$: INCL R4
 5B 54 E9 0453 1483 250\$: BLBC R4,260\$
 02 E0 0456 1484 BBS #PSLSV_Z,R11,270\$
 5B 08 88 045A 1485 BISB #PSLSM_N,R11
 ; Is next longword zero?
 ; No, leave loop
 ; Discard saved string descriptor
 ; Assume final result is positive
 ; Retrieve original R0/R1 pair
 ; Get byte count for first input string
 ; Point R1 to byte containing sign
 ; R0 contains the sign "digit"
 ; Dispatch on sign digit
 ; 10 => sign is '+'
 ; 11 => sign is '-'
 ; 12 => sign is '++'
 ; 13 => sign is '--'
 ; 14 => sign is '++'
 ; 15 => sign is '--'
 ; Count a minus sign
 ; Now check second input sign
 ; No real minus signs so far
 ; Retrieve original R2/R3 pair
 ; Get byte count for second input string
 ; Point R3 to byte containing sign
 ; R2 contains the sign "digit"
 ; Dispatch on sign digit
 ; 10 => sign is '+'
 ; 11 => sign is '-'
 ; 12 => sign is '++'
 ; 13 => sign is '--'
 ; 14 => sign is '++'
 ; 15 => sign is '--'
 ; Remember that sign was minus
 ; Even parity indicates positive result
 ; Step out of line for minus zero check
 ; Set N-bit in saved PSW

56 D6 045D 1486 255\$: INCL R6 ; Change sign to minus
045F 1487
045F 1488 MARK_POINT MULP_DIVP_0
045F 1489 260\$: INSV R6,#0,#4,(R7) ; Store sign in result string
67 0463
10 AE D4 0464 1490 CLRL 16(SP)
FB96' 31 0467 1491 BRW VAX\$DECIMAL_EXIT ; Set saved R4 to zero
046A 1492 ; Join common exit code
046A 1493 ; If the result is negative zero, then it must be changed to positive zero
046A 1494 ; unless overflow has occurred, in which case, the sign is left as negative
046A 1495 ; but the N-bit is clear.
046A 1496
EF 5B 01 E0 046A 1497 270\$: BBS #PSL\$V_V,R11,255\$; Make sign negative if overflow
EF 11 046E 1498 BRB 260\$; Sign will be positive

0470 1500 .SUBTITLE EXTEND_STRING_MULTIPLY - Multiply a String by a Number
 0470 1501 :+
 0470 1502 Functional Description:
 0470 1503 This routine multiplies an array of numbers (each array element LEQU 99) by a number (also LEQU 99). The resulting product array is added to another array, each of whose elements is also LEQU 99.
 0470 1504
 0470 1505
 0470 1506
 0470 1507
 0470 1508 Input Parameters:
 0470 1509
 0470 1510 R3 - Pointer to output array
 0470 1511 R4 - Input array size
 0470 1512 R5 - Input array address
 0470 1513 R6 - Multiplier
 0470 1514
 0470 1515 Output Parameters:
 0470 1516
 0470 1517 None
 0470 1518
 0470 1519 Implicit Output:
 0470 1520 The output array is altered.
 0470 1521 An intermediate product array is produced by multiplying each input array element by the multiplier. Each product array element is then added to the corresponding output array element.
 0470 1522
 0470 1523
 0470 1524
 0470 1525
 0470 1526
 0470 1527 Side Effects:
 0470 1528 R3, R4, and R5 are modified by this routine.
 0470 1529 R6 is preserved.
 0470 1530 R0, R1, and R2 are used as scratch registers. R0 and R1 contain the
 0470 1531 quadword result of EMUL that is then passed into EDIV.
 0470 1532
 0470 1533 Assumptions:
 0470 1534 This routine assumes that all array elements lie in the range from 0 to 99 inclusive. (This is true if all input strings contain only legal decimal digits.) The arithmetic performed by this routine will maintain this assumption. That is,
 0470 1535
 0470 1536
 0470 1537
 0470 1538 times input array element LEQU 99
 0470 1539 multiplier LEQU 99
 0470 1540 -----
 0470 1541 plus product LEQU 99*99
 0470 1542 carry
 0470 1543 -----
 0470 1544 plus modified product LEQU 99*100
 0470 1545 old output array element
 0470 1546 -----
 0470 1547 new output array element LEQU 99*101 = 9999
 0470 1548
 0470 1549
 0470 1550
 0470 1551
 0470 1552
 0470 1553
 0470 1554 A number LEQU 9999, when divided by 100, is guaranteed to produce both
 0470 1555 a quotient and a remainder LEQU 99.
 0470 1556 :-

			0470	1557				
			0470	1558	EXTEND_STRING_MULTIPLY:			
			0470	1559	CLRL R2			; Initial carry is zero
			0472	1560				
52	85	56	7A	0472	1561	10\$: EMUL R6,(R5)+,R2,R0		; Form modified product (R0 LEQU 9900)
		50		0476				
		50	63	C0	0477	1562 ADDL2 (R3),R0		; Add old output array element
00000064		8F	7B	047A	1563 EDIV #100,R0,R2,(R3)+			; Remainder to output array
83	52	50		0480				
				0483	1564			; Quotient becomes carry
	EC	54	F5	0483	1565 SOBGTR R4,10\$; Keep going?
				0486	1566			
				0486	1567	: This remaining code looks more complicated than it actually is. In the		
				0486	1568	: usual case, the routine exits immediately. In the event that a carry		
				0486	1569	: occurs, one additional entry in the output array will be modified. Only in		
				0486	1570	: the rare case of an output array consisting of a string of 99s will any		
				0486	1571	: significant looping occur.		
				0486	1572			
63	52	C0	0486	1573	ADDL2 R2,(R3)			; Add final carry
00000064	63	D1	0489	1574	20\$: CMPL (R3),#100			; Do we overflow into next digit pair?
			048B					
	01	1E	0490	1575	BGEQU 30\$; Branch if carry
		05	0492	1576	RSB			; Otherwise, all done
			0493	1577				
00000064	8F	C2	0493	1578	30\$: SUBL #100,(R3)+			; Readjust entry and advance pointer
83			0499					
63	D6	049A	1579		INCL (R3)			; Propogate carry
EB	11	049C	1580		BRB 20\$; ... and test this entry for overflow

049E 1582
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049E 1596
049E 1597
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049E 1599
049E 1600
049E 1601
049E 1602
049E 1603
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.SUBTITLE VAX\$DIVP - Divide Packed

Functional Description:

The dividend string specified by the dividend length and dividend address operands is divided by the divisor string specified by the divisor length and divisor address operands. The quotient string specified by the quotient length and quotient address operands is replaced by the result.

Input Parameters:

R0 - divrlen.rw	Number of digits in divisor string
R1 - divraddr.ab	Address of divisor string
R2 - divdlen.rw	Number of digits in dividend string
R3 - divdaddr.ab	Address of dividend string
R4 - quorlen.rw	Number of digits in quotient string
R5 - quoaddr.ab	Address of quotient string

Output Parameters:

R0 = 0
R1 = Address of the byte containing the most significant digit of the divisor string
R2 = 0
R3 = Address of the byte containing the most significant digit of the dividend string
R4 = 0
R5 = Address of the byte containing the most significant digit of the string containing the quotient

Condition Codes:

N <- quotient string LSS 0
Z <- quotient string EQL 0
V <- decimal overflow
C <- 0

Register Usage:

This routine uses all of the general registers. The condition codes are computed at the end of the instruction as the final result is stored in the quotient string. R11 is used to record the condition codes.

Algorithm:

This algorithm is the straightforward approach described in

The Art of Computer Programming
Second Edition

Volume 2 / Seminumerical Algorithms
Donald E. Knuth

1981
Addison-Wesley Publishing Company

049E 1639 : Reading, Massachusetts

049E 1640 : Notes:

049E 1641 : The choice of a longword array to store the auotient deserves a comment. In VAX\$MULP, a longword array was used because its elements were used directly by MULP and DIVP instructions. The use of longwords eliminated the need to convert back and forth between longwords and bytes. In this routine, the QUOTIENT DIGIT routine returns its result in a register, which result can easily be stored in whatever way is convenient. By using longwords instead of bytes, this routine can use the same end processing code as MULP, a sizeable savings in code.

049E 1642 :-

049E 1643 .ENABLE LOCAL_BLOCK

049E 1644 :+ This code path is entered if the divisor is zero.

049E 1645 : Input Parameter:

049E 1646 : (SP) - Return PC

049E 1647 : Output Parameters:

049E 1648 : 0(SP) - SRMSK_FLT_DIV T (Arithmetic trap code)

049E 1649 : 4(SP) - Final state PSL

049E 1650 : 8(SP) - Return PC

049E 1651 : Implicit Output:

049E 1652 : Control passes through this code to VAX\$REFLECT_TRAP.

049E 1653 :-

049E 1654 : DIVIDE_BY_ZERO:

049E 1655 : POPR #^M<R0,R1,R2,R3,R4,R5,R6,R7,R8,R9,R10,R11>

049E 1656 : Restore registers and reset SP

049E 1657 : MOVPSL -(SP)

049E 1658 : Save final PSL on stack

049E 1659 : PUSHL #SRMSK_FLT_DIV T

049E 1660 : Store arithmetic trap code

049E 1661 : BRW VAX\$REFLECT_TRAP

049E 1662 : Report exception

049E 1663 : If the divisor contains more nonzero digits than the dividend, then the

049E 1664 : quotient will be identically zero. Set up the stack and the registers (R4,

049E 1665 : R5, and R9) so that the exit code will be entered to produce this result.

049E 1666 : 04A9 1670 : CLRL -(SP)

049E 1667 : Count that digit

049E 1668 : MOVL #1,R9

049E 1669 : Store the zero in the output string

049E 1670 : BRW MULTIPLY_DIVIDE_EXIT

049E 1671 : 04B1 1684 : 1\$: CLRL -(SP)

049E 1672 : Save the lot

049E 1673 : 04B1 1685 : MOVL #1,R9

049E 1674 : ESTABLISH_HANDLER -

049E 1675 : ARITH_ACCVIO

049E 1676 : Store address of access

049E 1677 : violation handler

049E 1678 : 04B1 1688 : VAX\$DIVP::

049E 1679 : 04B1 1689 : PUSHR #^M<R0,R1,R2,R3,R4,R5,R6,R7,R8,R9,R10,R11>

049E 1680 : 04B1 1690 : ROPRAND_CHECK R4

049E 1681 : 04B1 1691 : Insure that R4 is LEQU 31

049E 1682 : 04B1 1692 :

049E 1683 : 04B1 1693 :

049E 1684 : 04B1 1694 :

049E 1685 : 04B1 1695 :

FB30' 30 04C5 1696 ROPRAND CHECK R2 ; Insure that R2 is LEQU 31
 04CD 1697 MARK_POINT DIVP BSBW 0
 04CD 1698 BSBW- DECIMALSSTRIP_ZEROS_R2_R3 ; Strip high order zeros from R2/R3
 04D0 1699
 04D0 1700 ROPRAND CHECK R0 ; Insure that R0 is LEQU 31
 04DB 1701 MARK_POINT DIVP BSBW 0
 FB25' 30 04DB 1702 BSBW- DECIMALSSTRIP_ZEROS_R0_R1 ; Strip high order zeros from R0/R1
 04DB 1703
 04DB 1704 : Insure that the divisor is not zero. Because leading zeros have already
 04DB 1705 : been eliminated, the divisor can only be zero if R0 is 0 (zero length
 04DB 1706 : strings are identically zero) or 1 (R1 contains a sign digit in the low
 04DB 1707 : order nibble and zero in the high order nibble). Note that an exception
 04DB 1708 : will not be generated if an even length string has an illegal nonzero digit
 04DB 1709 : stored in its most significant nibble (including an illegal form of a zero
 04DB 1710 : length string).

50 04 01 EF 04DB 1711 EXTZV #1,#4,R0,R0 ; Convert divisor digit count to bytes
 50 04DF
 06 12 04E0 1712 BNEQ 10\$; Skip zero divisor check unless zero
 61 F0 8F 93 04E2 1713 MARK_POINT DIVP 0
 B6 13 04E2 1714 BITB #^B11110000,7R1) ; Check for zero in ones digit
 04E6 1715 BEQL DIVIDE_BY_ZERO ; Generate exception if zero
 04EB 1716
 04EB 1717 : This routine chooses to do its work with a fair amount of internal storage,
 04EB 1718 : all of it allocated on the stack. The quotient is stored as it is computed,
 04EB 1719 : in a 16-longword array. The dividend and divisor are stored as longword arrays,
 04EB 1720 : with each array element storing a digit pair from the original packed
 04EB 1721 : decimal string. The numerator digits are shifted by one digit (multiplied
 04EB 1722 : by ten) so that the quotient has its digits correctly placed, leaving room
 04EB 1723 : for a sign in the low order nibble of the least significant byte. A scratch
 04EB 1724 : array is also allocated on the stack to accommodate intermediate results
 04EB 1725 : of the QUOTIENT_DIGIT routine.
 04EB 1726
 04EB 1727

58 50 D6 04E8 1728 10\$: INCL R0 ; Include least significant digit
 50 7D 04EA 1729 MOVQ R0,R8 ; Let R8 and R9 describe the divisor
 04ED 1730
 52 04 01 EF 04ED 1731 EXTZV #1,#4,R2,R2 ; Convert dividend digit count to bytes
 52 04F1
 7E 52 D6 04F2 1732 INCL R2 ; Include least significant digit
 52 7D 04F4 1733 MOVQ R2,-(SP) ; Save dividend descriptor on stack
 04F7 1734
 56 52 50 C3 04F7 1735 SUBL3 R0,R2,R6 ; Calculate main loop count
 AC 1F 04FB 1736 BLSSU 1\$; Quotient will be zero
 56 D6 04FD 1737 INCL R6 ; One extra digit is always there
 04FF 1738
 04FF 1739 : Allocate R6 longwords of zero on the stack
 04FF 1740

50 56 D0 04FF 1741 MOVL R6,R0 ; Let R0 be the loop counter
 7E D4 0502 1742 15\$: CLRL -(SP) ; Set aside another quotient digit
 FB 50 F5 0504 1743 SOBGTR R0,15\$; Keep going
 0507 1744
 57 5E D0 0507 1745 MOVL SP,R7 ; Remember where this array starts
 050A 1746

050A 1747 : The divisor will be stored on the stack as an array of
 050A 1748 : longwords. Each array element contains a number between 0 and 99,
 050A 1749 : representing a pair of digits in the original packed decimal string.
 050A 1750 : Because the units digit is stored with the sign in packed decimal format.

050A 1751 : it is necessary to shift the number as we store it. This is accomplished by
 050A 1752 : multiplying the number by ten.
 050A 1753
 050A 1754 : The divisor string is described by R8 (byte count) and R9 (address of most
 050A 1755 : significant digit pair).
 050A 1756

55 58 59 C1 050A 1757 ADDL3 R9,R8,R5 ; Point R5 beyond sign digit
 54 58 D0 050E 1758 MOVL R8,R4 ; R4 contains the loop count
 0511 1759
 0511 1760 : Put in an extra digit place for the divisor. This allows several common
 0511 1761 : subroutines to be used when operating on the divisor string.
 0511 1762
 7E D4 0511 1763 CLRL -(SP) ; Set aside a place holder
 0513 1764
 0513 1765 : An array of longwords is allocated on the stack. R3 starts out pointing
 0513 1766 : at the longword beyond the top of the stack. The first remainder, guaranteed
 0513 1767 : to be zero, is "stored" here. The rest of the digit pairs are stored safely
 0513 1768 : below the top of the stack.
 0513 1769

53 5E 53 58 CE 0513 1770 MNEGL R8,R3 ; Stack grows toward lower addresses
 53 5E 6E43 DE 0516 1771 MOVAL (SP)[R3],SP ; Allocate the space
 04 C3 051A 1772 SUBL3 #4,SP,R3 ; Point R3 at next lower longword
 051E 1773
 051E 1774 MARK POINT DIVP_R6_R7
 51 51 75 9A 051E 1775 20\$: MOVZBL -(R5),R1 ; Get next digit pair
 0000'CF41 9A 0521 1776 MOVZBL DECIMAL\$PACKED_TO_BINARY_TABLE[R1],-
 52 51 0A 7A 0527 1777 R1 ; Convert digits to binary
 50 052B
 00000064 8F 7B 052C 1778 EMUL #10,R1,R2,R0 ; Multiply by 10
 83 52 50 0532 1779 EDIV #100,R0,R2,(R3)+ ; Divide by 100
 E6 54 F5 0535 1780 SOBGTR R4,20\$
 0538 1781
 0538 1782 : There are two cases where the final quotient (contents of R2) is zero.
 0538 1783 : In these cases, the number of nonzero digit pairs in the divisor array is
 0538 1784 : smaller by one than the number of bytes containing the original packed decimal
 0538 1785 : string. One case is a divisor string with an even number of digits. The
 0538 1786 : second case is a divisor string with an odd number of digits but the most
 0538 1787 : significant digit is zero (essentially a variant of the first case). The
 0538 1788 : simplest way to handle all of these cases is to decrement R8, the divisor
 0538 1789 : counter, if R2 is zero. Note that previous checks for a zero divisor
 0538 1790 : prevent R8 from going to zero.
 0538 1791

63 52 D0 0538 1792 MOVL R2,(R3) ; Store final quotient
 0A 12 053B 1793 BNEQ 25\$; Leave well enough alone if nonzero
 56 D6 053D 1794 INCL R6 ; One more quotient digit
 57 04 C2 053F 1795 SUBL #4,R7 ; Make room for it
 58 D7 0542 1796 DECL R8 ; Count one less divisor "digit"
 01 12 0544 1797 BNEQ 25\$

0546 1798 : ***** BEGIN TEMP *****
 0546 1800
 0546 1801 : THE FOLLOWING HALT INSTRUCTION SHOULD BE REPLACED WITH THE CORRECT
 0546 1802 : ABORT CODE.
 0546 1803
 0546 1804 : THE HALT IS SIMILAR TO THE
 0546 1805

```

      0546 1806 ;:: MICROCODE CANNOT GET HERE
      0546 1807 ;:: ERRORS THAT OTHER IMPLEMENTATIONS USE.
      0546 1808 ;:: halt ; This will cause an OPCDEC exception
      00 0546 1810 ;:: ***** END TEMP *****
      0547 1811 ;:: *****
      0547 1812 ;:: *****
      0547 1813 ;:: *****
      59 SE DO 0547 1814 25$: MOVL SP,R9 ; R9 locates low order divisor digit
      054A 1815
      054A 1816 ; The dividend is stored on the stack as an array of longwords. It does not
      054A 1817 ; have its digit pairs shifted so that this storage loop is simpler. An extra
      054A 1818 ; place is set aside in the event that it is necessary to normalize the
      054A 1819 ; dividend and divisor before division is attempted.
      054A 1820
      52 7E D4 054A 1821 CLRL -(SP) ; Set aside space for U[0]
      52 6746 DE 054C 1822 MOVAL (R7)[R6],R2 ; Retrieve dividend descriptor
      52 62 7D 0550 1823 MOVQ (R2),R2 ; ... in two steps
      0553 1824
      7E 51 83 9A 0553 1825 MARK POINT DIVP_R6_R7
      0000'CF41 9A 0556 1826 30$: MOVZBL (R3)+,R1 ; Get next decimal digit pair
      F4 52 F5 055C 1827 MOVZBL DECIMALSPACKED_TO_BINARY TABLE[R1],-
      055F 1828 -(SP) ; Convert digits to binary
      055F 1829 SOBGTR R2,30$ ; Loop through entire input string
      055F 1830
      055F 1831 ; From this point until the common exit path for MULP and DIVP is entered,
      055F 1832 ; no access violations that need to be backed out can occur. We do not need
      055F 1833 ; to keep the address of ARITH_ACCVIO in R10 for this stretch of code. Note
      055F 1834 ; that R10 must be reloaded before the exit code executes because the
      055F 1835 ; destination string is written and may cause access violations.
      055F 1836
      5A 6746 DO 055F 1837 MOVL (R7)[R6],R10 ; Retrieve size of dividend array
      5B SE DO 0563 1838 MOVL SP,R11 ; R11 locates low order dividend digit
      0566 1839
      0566 1840 ; Allocate a scratch array on the stack the same size as the divisor array
      0566 1841 ; (which is one larger than the number of digit pairs)
      0566 1842
      SE 52 58 CE 0566 1843 MNEGL R8,R2 ; Need a negative index
      FC AE42 DE 0569 1844 MOVAL -4(SP)[R2],SP ; Adjust stack pointer
      056E 1845
      056E 1846 +
      056E 1847 ; At this point, the stack and relevant general registers contain the
      056E 1848 ; following information. In this description, N represents the number
      056E 1849 ; of digit pairs in the divisor and M represents the number of digit
      056E 1850 ; pairs in the dividend.
      056E 1851
      056E 1852 scratch +-----+ <- SP
      056E 1853 | N+1 longwords |
      056E 1854 +-----+ <- R11
      056E 1855 dividend +-----+
      056E 1856 | M+1 longwords |
      056E 1857 +-----+ <- R9
      056E 1858 divisor +-----+
      056E 1859 quotient +-----+ <- R7
      056E 1860 | M+1-N longwords |
      056E 1861 +-----+ <- R0..R11
      056E 1862
  
```

056E 1863 ;
 056E 1864 ; R6 - Number of longwords in quotient array (M+1-N)
 056E 1865 ; R7 - Address of beginning of quotient array
 056E 1866 ; R8 - Number of digit pairs in divisor (called N)
 056E 1867 ; R9 - Address of low order digits in divisor
 056E 1868 ; R10 - Number of digit pairs in dividend (called M)
 056E 1869 ; R11 - Address of low order digits in dividend
 056E 1870 ;-

7E 6E DF 056E 1871
 7E 5B 7D 0570 1872 PUSHAL (SP) ; Store address of scratch array
 7E 5A 7D 0573 1873 MOVQ R8,-(SP) ; Remember divisor descriptor
 0576 1874 MOVQ R10,-(SP) ; Remember dividend descriptor

'576 1875 : The algorithm that guesses the quotient digit can be guaranteed to be off
 0576 1876 : by no more than two if the high order digit of the divisor (called V[1]) is
 0576 1877 : at least as large as 50 (our radix divided by 2). If the high order digit
 0576 1878 : is too small, we "normalize" the numerator and denominator by multiplying
 0576 1879 : them by the same number, namely 100/(V[1]+1).
 0576 1880 :

FC A948 01 C1 0576 1882 ADDL3 #1,-4(R9)[R8],R0 ; Compute V[1] + 1
 50 057B
 33 50 D1 057C 1883 CMPL R0,#51 ; Compare to 50 + 1
 14 18 057F 1884 BGEQ 40\$; Skip normalization if V[1] big enough
 50 C7 0581 1885 DIVL3 R0,#100,R3 ; Compute normalization factor

00000064 8F 0583
 53 0588
 54 58 7D 0589 1886 MOVQ R8,R4 ; Get descriptor of divisor
 00E0 30 058C 1887 BSBW MULTIPLY_STRING ; Normalize divisor
 54 5A 7D 058F 1888 MOVQ R10,R4 ; Get descriptor of dividend
 00DA 30 0592 1889 BSBW MULTIPLY_STRING ; Normalize dividend

0595 1890 : We have now reached the point where we can start calculating quotient digits.
 0595 1891 : In the following loop, R5 and R6 are loop invariants. R5 contains the number
 0595 1892 : of digit pairs in the divisor. R6 always points to the longword beyond the
 0595 1893 : most significant digit in the dividend string. R7 and R8 must be loaded on
 0595 1894 : each pass through because these two pointers are modified. Notice that the
 0595 1895 : address of the divisor array is exactly what we want to store in R6.
 0595 1896 :

5A 56 7D 0595 1898 40\$: MOVQ R6,R10 ; Let R10/R11 describe quotient and loop
 5B DD 0598 1899 PUSHL R11 ; Save quotient address for exit code
 5B 6B4A DE 059A 1900 MOVAL (R11)[R10],R11 ; Store quotient digits from high end

059E 1901 : This rather harmless looking loop is where the work is done
 059E 1902 :

55 58 7D 059E 1904 MOVQ R8,R5 ; Initialize count and dividend address
 59 5A D0 05A1 1905 MOVL R10,R9 ; Remember the loop count in R9

57 10 AE 7D 05A4 1907 50\$: MOVQ 16(SP),R7 ; Load divisor and scratch addresses
 001F 30 05A8 1908 BSBW QUOTIENT_DIGIT ; Get the next quotient digit
 7B 53 D0 05AB 1909 MOVL R3,-(R11) ; Store it
 56 04 C2 05AE 1910 SUBL #4,R6 ; "Advance" dividend pointer
 F0 5A F5 05B1 1911 SOBGTR R10,50\$; ... and go back for more

05B4 1912 : The quotient digits have been stored on the stack. Eliminate the rest of the
 05B4 1913 : stack storage and enter the completion code that this routine shares with
 05B4 1914 : VAX\$MULP. Note that R9 is already set up with the longword count used by
 05B4 1915 : the exit code. Note also that R11 is pointing to the saved dividend descriptor
 05B4 1916 :

05B4 1917 ; that sits on top of the saved register array.
05B4 1918
54 5E 6E D0 05B4 1919 MOVL (SP),SP : Reset stack pointer
18 AB49 DE 05B7 1920 MOVAL <<4+2>>+- : Skip over saved dividend descriptor
54 64 7D 05BC 1921 <<4+4>>(R11)[R9],R4 ; and retrieve original R4 and R5
05BF 1922 MOVQ (R4),R4 ; ... in two steps
05BF 1923
05BF 1924 ; The following is a HACK.
05BF 1925
05BF 1926 ; The method used to obtain quotient digits generally leaves garbage (nonzero)
05BF 1927 ; in what will become the sign digit. (In fact, this is the tenths digit of a
05BF 1928 ; decimal expansion of the remainder.) We need to make the least significant
05BF 1929 ; digit a multiple of ten.
05BF 1930
50 6E 0A C7 05BF 1931 DIVL3 #10,(SP),R0 ; Divide by ten, losing remainder
6E 50 0A C5 05C3 1932 MULL3 #10,R0,(SP) ; Store only tens digit
FDB3 31 05C7 1933 BRW MULTIPLY_DIVIDE_EXIT ; Join common exit code
05CA 1934
05CA 1935
.DISABLE LOCAL_BLOCK

05CA 1938
05CA 1939
05CA 1940
05CA 1941
05CA 1942
05CA 1943
05CA 1944
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05CA 1946
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05CA 1948
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05CA 1991
05CA 1992
05CA 1993
05CA 1994

.SUBTITLE QUOTIENT_DIGIT - Get Next Digit in Quotient

Functional Description:

This routine divides an (N+1)-element array of longwords by an N-element array, producing a single quotient digit in the range of 0 to 99 inclusive. The dividend array is modified by subtracting the product of the divisor array and the quotient digit.

The "numbers" that this array operates on multiple precision numbers in radix 100. Each digit (a number between 0 and 99) is stored in a longword array element with more significant digits stored at higher addresses. The dividend string and the scratch string (also called the product string) contain one more element than the divisor string.

Input Parameters:

R5 - Number of "digits" (array elements) in divisor array (preserved)
R6 - Address of longword immediately following most significant digit of dividend string (preserved)
R7 - Address of least significant digit in divisor string (modified)
R8 - Address of least significant digit in product string (modified)

Output Parameters:

R3 - The quotient that results from dividing the dividend string by the divisor string.

The final states of the three pointer registers are listed here for completeness.

R6 - Address of longword immediately following most significant digit of dividend string

R7 - Address of longword immediately following most significant digit of divisor string. This longword must always contain zero.

R8 - Address of longword immediately following most significant digit of product string

Implicit Output:

The contents of the dividend array are modified to reflect the subtraction of the product string. The result of this subtraction could be stored elsewhere. It is a convenience to store it in the dividend array on top of those array elements that are no longer needed.

The contents of the divisor array are preserved.

Side Effects:

R7 and R8 are modified by this routine. (See implicit output list.)

R5 and R6 are preserved.

R0, R1, R2, and R4 are used as scratch registers. R0 and R1 contain the

05CA 1995 : quadword result of EMUL that is then passed into EDIV. R2 is the
 05CA 1996 : carry from one step to the next. R4 is the loop counter.
 05CA 1997 :-
 05CA 1998

00000064 8F FB A6 FC A6 50 50	7A 05CA 2000 QUOTIENT DIGIT: 05D0 05D4	05CA 2001 DIVL2 -4(R7)[R5],R0 ; R0 <- R0 / V[1] 05DA 2002 MOVL R0,R3 05DD 2003 BEQL 45\$; Store quotient "digit" in R3 05DF 2004 CMPL R3,#100 ; Nothing to do if quotient is zero 05E1 2005 BLSSU 5\$; Is quotient LEQU 99? 05E8 2006 MOVL #99,R3 ; Branch if quotient OK 05EE 2007 ; Otherwise start with 99	
		05EF 2008 : We will now multiply the divisor array by the quotient digit, storing the 05EF 2009 ; product in the scratch array.	
		05EF 2010 CLRL R2 ; Start out with a carry of zero 05F1 2012 MOVL R5,R4 ; R4 will be the loop counter	
		05F4 2013 05F8 2014 10\$: EMUL R3,(R7)+,R2,R0 ; Multiply next divisor digit	
		05FF 2015 EDIV #100,R0,R2,(R8)+ ; Remainder to input array	
		0602 2016 SOBGTR R4,10\$; Quotient becomes carry 0605 2017 0605 2018 MOVL R2,(R8)+ ; More divisor digits?	
		0605 2019 0608 2020 MOVL R2,(R8)+ ; Store final carry	
		0608 2021 : If the product array is larger than the dividend array, then the quotient is 0608 2022 ; too large. To avoid a second trip through the rather costly EMUL/EDIV loop, 0608 2023 ; and also to avoid array subtraction that produces a negative result, we will 0608 2024 ; first compare the product and dividend arrays. If the product is smaller, we 0608 2025 ; can safely subtract. If the product is larger, we decrease the quotient by 0608 2026 ; one and subtract the divisor array from the product array.	
		0608 2027 0608 2028 15\$: MOVL R6,R0 ; Point R0 and R1 to high address ends 0608 2029 MOVL R8,R1 ; of dividend and scratch strings	
		060E 2030 MOVL R5,R4 ; Initialize the loop counter	
		0611 2031 0611 2032 ; The comparison is done from most to least significant digits 0611 2033	
		0611 2034 20\$: CMPL -(R1),-(R0) ; Compare next pair of digits 0611 2035 BLSSU 30\$; Leave loop if product is smaller	
		0611 2036 BGTRU 50\$; Also leave if product is larger	
		0618 2037 SOBGEQ R4,20\$; More to test?	
		0618 2038 0618 2039 : If we drop through the loop, then the dividend and product are equal. We 0618 2040 ; simply store zeros in the dividend array (the equivalent of subtraction 0618 2041 ; of equal arrays) and return. Note that R0 is already pointing to the 0618 2042 ; least significant dividend array element.	
		0618 2043 061B 2044 MOVL R5,R4 ; Initialize still another loop counter	
		061E 2045	

80	D4	061E	2046	25\$: CLRL (R0)+ FB 54 F4 0620 2047 SOBGEQ R4,25\$; Store another zero	; Keep going?
			05	0623 2048 ; RSB ; Return to caller	
			0624 2050		
			0624 2051 : If we drop through the loop, then the quotient that is stored in R3 is good.		
			0624 2052 : We need to subtract the product array from the dividend array. Note that R0		
			0624 2053 : and R1 need to be adjusted to point to the least significant array elements		
			0624 2054 : before the subtraction can begin.		
			0624 2055		
50	54	CE	0624 2056 30\$: MNEGL R4,R4	; We need a negative index	
51	6044	DE	0627 2057 MOVAL (R0)[R4],R0	; Adjust dividend pointer	
51	6144	DE	062B 2058 MOVAL (R1)[R4],R1	; ... and product pointer	
54	55	DO	062F 2059 MOVL R5,R4	; R4 will count still another loop	
80	81	C2	0632 2060 35\$: SUBL2 (R1)+,(R0)+	; Subtract next digits	
00000064	8F	CO	0635 2062 BGEQ 40\$; Skip to end of loop if no borrow	
FC	A0		0637 2063 ADDL2 #100,-4(R0)	; Add borrow back to this digit	
60	D7	063D	2064 DECL (R0)		
EE 54	F4	0641	2065 40\$: SOBGEQ R4,35\$; ... and borrow from next highest digit	
		0644	2066	; Keep going?	
		0644	2067 : This is the exit path. R3 contains the quotient digit. The pointers to the		
		0644	2068 : various input and scratch arrays are in an indeterminate state.		
		0644	2069		
		05	0644 2070 45\$: RSB	; Return to caller	
		0645	2071		
		0645	2072 : The first guess at the quotient digit is too large. The brute force		
		0645	2073 : approach is to decrement the quotient by one and execute the EMUL/EDIV loop		
		0645	2074 : again. Note, however, that we can evaluate the modified product by		
		0645	2075 : subtracting the divisor from the initial product. Note also that, because		
		0645	2076 : the leading digit in the divisor is "large enough", we can only end up in		
		0645	2077 : this code path twice. (That is, the initial guess at the quotient will		
		0645	2078 : never be off by more than two.)		
53	D7	0645 2080 50\$: DECL R3	; Try quotient smaller by one		
FB	13	0647 2081 BEQL 45\$; All done if zero		
		0649 2082			
		0649 2083 : Point R1 and R2 at the least significant digits of the scratch and product			
		0649 2084 : strings respectively.			
		0649 2085			
51	50	CE	0649 2086 MNEGL R5,R0	; Need a negative index	
51	FC	DE	064C 2087 MOVAL -4(R8)[R0],R1	; Scratch array contains N+1 elements	
52	6740	DE	0651 2088 MOVAL (R7)[R0],R2	; Product array contains N elements	
54	55	DO	0655 2089 MOVL R5,R4	; R4 will count still another loop	
81	82	C2	0658 2091 60\$: SUBL2 (R2)+,(R1)+	; Subtract next digits	
00000064	8F	CO	065B 2092 BGEQ 70\$; Skip to end of loop if no borrow	
FC	A1		065D 2093 ADDL2 #100,-4(R1)	; Add borrow back to this digit	
61	D7	0663	2094 DECL (R1)		
EE 54	F4	0665 2095 70\$: SOBGEQ R4,60\$; ... and borrow from next highest digit		
51	04	CO	066A 2096 ADDL2 #4,R1	; Keep going?	
99	11	066D 2098 BRB 15\$; Point R1 at most significant digit		

066F	2100		.SUBTITLE	MULTIPLY_STRING - Multiply a String by a Number	
066F	2101		+ Functional Description:		
066F	2102			This routine multiplies an array of numbers (each array element LEQU 99) by a number (also LEQU 99). Each array element in the input array is replaced with the modified product, with the carry propagated to the next array element.	
066F	2103				
066F	2104				
066F	2105				
066F	2106				
066F	2107				
066F	2108				
066F	2109				
066F	2110				
066F	2111				
066F	2112				
066F	2113				
066F	2114				
066F	2115				
066F	2116				
066F	2117				
066F	2118				
066F	2119				
066F	2120				
066F	2121				
066F	2122				
066F	2123				
066F	2124				
066F	2125				
066F	2126				
066F	2127				
066F	2128				
066F	2129				
066F	2130				
066F	2131				
066F	2132				
066F	2133				
066F	2134				
066F	2135				
066F	2136				
066F	2137				
066F	2138				
066F	2139				
066F	2140				
066F	2141				
066F	2142				
066F	2143	MULTIPLY STRING:			
52	D4	066F	2144	CLRL R2	: Initial carry is zero
		066F	2145	0671	
52	65	53	7A	0671	2146 10\$: EMUL R3,(R5),R2,R0 ; Form modified product (R0 LEQU 9900)
				0675	
00000064	8F	7B	0676	2147	EDIV #100,R0,R2,(R5)+ ; Remainder to input array
85	52	50		067C	
				067F	2148
				067F	
EF	54	F5	0682	2151	SOBGTR R4,10\$; Quotient becomes carry
			0682	2150	;
65	52	D0	0682	2151	MOVL R2,(R5) ; Keep going?
			0685	2152	RSB ; Store final carry

```

0686 2154 .SUBTITLE      DECIMAL_ROPRAND
0686 2155 :- Functional Description:
0686 2156
0686 2157
0686 2158 This routine receives control when a digit count larger than 31
0686 2159 is detected. The exception is architecturally defined as an
0686 2160 abort so there is no need to store intermediate state. All of the
0686 2161 routines in this module save all registers R0 through R11 before
0686 2162 performing the digit check. These registers must be restored
0686 2163 before control is passed to VAX$ROPRAND.
0686 2164
0686 2165 Input Parameters:
0686 2166
0686 2167 00(SP) - Saved R0
0686 2168 .
0686 2169
0686 2170 44(SP) - Saved R11
0686 2171 48(SP) - Return PC from VAX$xxxxxx routine
0686 2172
0686 2173 Output Parameters:
0686 2174
0686 2175 00(SP) - Offset in packed register array to delta PC byte
0686 2176 04(SP) - Return PC from VAX$xxxxxx routine
0686 2177
0686 2178 Implicit Output:
0686 2179
0686 2180 This routine passes control to VAX$ROPRAND where further
0686 2181 exception processing takes place.
0686 2182 :- DE
0686 2183
0686 2184 ASSUME ADDP6_B_DELTA_PC EQ ADDP4_B_DELTA_PC
0686 2185 ASSUME SUBP4_B_DELTA_PC EQ ADDP4_B_DELTA_PC
0686 2186 ASSUME SUBP6_B_DELTA_PC EQ ADDP4_B_DELTA_PC
0686 2187 ASSUME MULP_B_DELTA_PC EQ ADDP4_B_DELTA_PC
0686 2188 ASSUME DIVP_B_DELTA_PC EQ ADDP4_B_DELTA_PC
0686 2189
0686 2190 DECIMAL_ROPRAND:
0FFF 8F BA 0686 2191 POPR #^M<R0,R1,R2,R3,R4,R5,R6,R7,R8,R9,R10,R11>
03 DD 068A 2192 PUSHL #ADDP4_B_DELTA_PC ; Store offset to delta PC byte
F971' 31 068C 2193 BRW VAX$ROPRAND ; Pass control along

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068F 2195 .SUBTITLE ARITH_ACCVIO - Reflect an Access Violation
 068F 2196 :+ Functional Description:
 068F 2197 This routine receives control when an access violation occurs while
 068F 2198 executing within the emulator routines for ADDP4, ADDP6, SUBP4, SUBP6,
 068F 2199 MULP, or DIVP.
 068F 2200 The routine header for ASHP_ACCVIO in module VAX\$ASHP contains a
 068F 2201 detailed description of access violation handling for the decimal
 068F 2202 string instructions.
 068F 2203 Input Parameters:
 068F 2204 See routine ASHP_ACCVIO in module VAX\$ASHP
 068F 2205 Output Parameters:
 068F 2206 See routine ASHP_ACCVIO in module VAX\$ASHP
 068F 2207
 068F 2208
 068F 2209
 068F 2210
 068F 2211
 068F 2212
 068F 2213
 068F 2214 ;-
 068F 2215
 068F 2216 ARITH_ACCVIO:
 F96B 52 D4 068F 2217 CLRL R2 ; Initialize the counter
 CF 9F 0691 2218 PUSHAB MODULE_BASE ; Store base address of this module
 06FC'CF 9F 0695 2219 PUSHAB MODULE_END ; Store module end address
 F964' 30 0699 2220 BSBW DECIMA[\$BOUNDS_CHECK] ; Check if PC is inside the module
 SE 04 C0 069C 2221 ADDL #4,SP ; Discard end address
 51 8E C2 069F 2222 SUBL2 (SP)+,R1 ; Get PC relative to this base
 06A2 2223
 0000'CF42 51 B1 06A2 2224 10\$: CMPW R1,PC_TABLE_BASE[R2] ; Is this the right PC?
 07 13 06A8 2225 BEQL 30\$; Exit loop if true
 F4 52 2A F2 06AA 2226 AOBLS #TABLE_SIZE,R2,10\$; Do the entire table
 06AE 2227
 06AE 2228 ; If we drop through the dispatching based on PC, then the exception is not
 06AE 2229 ; one that we want to back up. We simply reflect the exception to the user.
 06AE 2230
 OF BA 06AE 2231 20\$: POPR #^M<R0,R1,R2,R3> ; Restore saved registers
 05 06B0 2232 RSB ; Return to exception dispatcher
 06B1 2233
 06B1 2234 ; The exception PC matched one of the entries in our PC table. R2 contains
 06B1 2235 ; the index into both the PC table and the handler table. R1 has served
 06B1 2236 ; its purpose and can be used as a scratch register.
 06B1 2237
 51 0000'CF42 3C 06B1 2238 30\$: MOVZWL HANDLER_TABLE BASE[R2],R1 ; Get the offset to the handler
 F944 CF41 17 06B7 2239 JMP MODULE_BASE[RT] ; Pass control to the handler
 06BC 2240
 06BC 2241 ; In all of the instruction-specific routines, the state of the stack
 06BC 2242 ; will be shown as it was when the exception occurred. All offsets will
 06BC 2243 ; be pictured relative to R0.

06BC 2245 .SUBTITLE Access Violation Handling for ADDPx and SUBPx
 06BC 2246 +
 06BC 2247 Functional Description:
 06BC 2248 The only difference among the various entry points is the number of
 06BC 2249 longwords on the stack. R0 is advanced beyond these longwords to point
 06BC 2250 to the list of saved registers. These registers are then restored,
 06BC 2251 effectively backing the routine up to its initial state.
 06BC 2252
 06BC 2253 Input Parameters:
 06BC 2254 R0 - Address of top of stack when access violation occurred
 06BC 2255 See specific entry points for details
 06BC 2256 Output Parameters:
 06BC 2257 See input parameter list for VAX\$DECIMAL_ACCVIO in module VAX\$ASHP
 06BC 2258
 06BC 2259
 06BC 2260
 06BC 2261
 06BC 2262 ADD_SUB_BSBW_24
 06BC 2263 -
 06BC 2264
 06BC 2265 +
 06BC 2266 An access violation occurred in one of the subroutines ADD_PACKED_BYTE,
 06BC 2267 SUB_PACKED_BYTE, or STORE_RESULT. In addition to the six longwords of work
 06BC 2268 space, this routine has an additional longword, the return PC, on the
 06BC 2269 stack.
 06BC 2270 00(R0) - Return PC in mainline VAX\$xxxxxx routine
 06BC 2271 04(R0) - Address of sign byte of destination string
 06BC 2272 08(R0) - First longword of scratch space
 06BC 2273 etc.
 06BC 2274 :-
 06BC 2275 ADD_SUB_BSBW_24:
 50 04 C0 06BC 2276 ADDL #4,R0 ; Skip over return PC and drop into ...
 06BF 2277
 06BF 2278
 06BF 2279 ADD_SUB_BSBW_24:
 06BF 2280 ADDL #4,R0 ; Skip over return PC and drop into ...
 06BF 2281
 06BF 2282 +
 06BF 2283 ADD_SUB_24
 06BF 2284
 06BF 2285 There are five longwords of workspace and a saved string address on the stack
 06BF 2286 for this entry point.
 06BF 2287 00(R0) - Address of sign byte of destination string
 06BF 2288 04(R0) - First longword of scratch space
 06BF 2289
 06BF 2290
 06BF 2291
 06BF 2292 20(R0) - Fifth longword of scratch space
 06BF 2293 24(SP) - Saved R0
 06BF 2294 28(SP) - Saved R1
 06BF 2295 etc.
 06BF 2296 :-
 06BF 2297
 06BF 2298 ADD_SUB_24:
 50 F93B' 18 C0 06BF 2299 ADDL #24,R0 ; Discard scratch space on stack
 06C2 2300 BRW VAX\$DECIMAL_ACCVIO ; Join common code to restore registers
 06C5 2301

06CE 2319 .SUBTITLE Access Violation Handling for MULP and DIVP
 06CE 2320 +
 06CE 2321 Functional Description:
 06CE 2322 The only difference among the various entry points is the number of
 06CE 2323 longwords on the stack. R0 is advanced beyond these longwords to point
 06CE 2324 to the list of saved registers. These registers are then restored,
 06CE 2325 effectively backing the routine up to its initial state.
 06CE 2326
 06CE 2327 Input Parameters:
 06CE 2328
 06CE 2329
 06CE 2330 R0 - Address of top of stack when access violation occurred
 06CE 2331
 06CE 2332 See specific entry points for details
 06CE 2333
 06CE 2334 Output Parameters:
 06CE 2335
 06CE 2336 See input parameter list for VAX\$DECIMAL_ACCVIO in module VAXSASHP
 06CE 2337 -
 06CE 2338
 06CE 2339 + MULP_R8
 06CE 2340
 06CE 2341 An access violation occurred while MULP was accessing one of its two source
 06CE 2342 strings. In this particular case, MULP was storing the longer of the two
 06CE 2343 input strings in a longword array on the top of the stack. There is an
 06CE 2344 array of R8 longwords on top of an array of 32 longwords on top of the
 06CE 2345 saved register array.
 06CE 2346
 06CE 2347 R8 - Number of longwords on top of the 32-longword array
 06CE 2348
 06CE 2349 -
 06CE 2350
 06CE 2351 .ENABLE LOCAL_BLOCK
 06CE 2352
 06CE 2353 MULP_R8:
 50 6048 DE 06CE 2354 MOVAL (R0)[R8],R0 ; Discard input array storage
 04 11 06D2 2355 BRB 10\$; Might as well share a little code
 06D4 2356
 06D4 2357 + MULP_AT_SP
 06D4 2358
 06D4 2359 An access violation occurred while MULP was accessing one of its two source
 06D4 2360 strings. In this case, the access violation occurred in the middle of the
 06D4 2361 grand multiply loop as a digit pair was being retrieved from the shorter of
 06D4 2362 the two input strings. The address of the start of the 32-longword array
 06D4 2363 was itself stored on top of the stack for convenience.
 06D4 2364
 06D4 2365 00(R0) - Saved byte count of longer input string
 06D4 2366 04(R0) - Saved address of longer input string
 06D4 2367 08(R0) - Address of 32-longword array farther down the stack
 06D4 2368
 06D4 2369 -
 06D4 2370
 06D4 2371 MULP_AT_SP:
 50 08 A0 D0 06D4 2372 MOVL 8(R0),R0 ; Locate start of 32-longword array
 F920' 31 06DB 2373 10\$: MOVAB <<4*32> + <4*2>>(R0),R0 ; Throw that away, too
 06DD 2374 BRW VAX\$DECIMAL_ACCVIO ; Join common code to restore registers
 06E0 2375

06E0 2376 .DISABLE LOCAL_BLOCK

06E0 2377

06E0 2378 :+ MULP_DIVP_R9

06E0 2381 An access violation occurred while the final result was being stored in the
06E0 2382 result string. In this common exit code path, R9 counts the number of
06E0 2383 longwords on the stack. In all cases where an access violation can occur, a
06E0 2384 longword has been removed from the stack but R9 has not yet been
06E0 2385 decremented to reflect this. The conceptual instruction sequence that
06E0 2386 resets the stack pointer (really R0) to point to the start of the saved
06E0 2387 register array is

06E0 2388

06E0 2389 DECL R9

06E0 2390 MOVAL (R0)[R9]

06E0 2391 A single instruction accomplishes this.

06E0 2392

06E0 2393

06E0 2394 R9 - One more than the number of longwords on the stack on top
06E0 2395 of the saved register array.

06E0 2396 00(R0) - First longword of scratch storage remaining on the stack

06E0 2397 .

06E0 2398 .

06E0 2399 .

06E0 2400 zz-4(R0) - Last longword of scratch storage

06E0 2401 zz+0(R0) - Saved count of dividend or multiplier string

06E0 2402 zz+4(R0) - Saved address of dividend or multiplier string

06E0 2403 zz+8(R0) - Saved R0

06E0 2404 zz+12(R0) - Saved R1

06E0 2405 etc.

06E0 2406

06E0 2407 where zz = 4 * (R9 - 1)

06E0 2408 :-

06E0 2409

50 04 A049, DE 06E0 2410 MULP_DIVP R9:
F918' 31 06E0 2411 MOVAL 4(R0)[R9],R0 ; Discard scratch storage on stack
06E5 2412 BRW VAX\$DECIMAL_ACCVIO ; Join common code to restore registers

06E8 2413

06E8 2414 :+ MULP_DIVP_8

06E8 2415

06E8 2416 An access violation occurred in the common exit path after the scratch array
06E8 2417 had been removed from the stack but before the saved descriptor for the
06E8 2418 multiplier string was discarded.

06E8 2419

06E8 2420 0(R0) - Saved count of dividend or multiplier string

06E8 2421 4(R0) - Saved address of dividend or multiplier string

06E8 2422 8(R0) - Saved R0

06E8 2423 12(R0) - Saved R1

06E8 2424 etc.

06E8 2425 :-

06E8 2426

06E8 2427

50 08 C0 06E8 2428 MULP_DIVP 8:
F912' 31 06E8 2429 ADDL #8,R0 ; Discard multiplier string descriptor
06EB 2430 BRW VAX\$DECIMAL_ACCVIO ; Join common code to restore registers

06EE 2431

06EE 2432 ;+

```

06EE 2433 : MULP_BSBW_0
06EE 2434 : DIVP_BSBW_0
06EE 2435
06EE 2436 : An access violation occurred in one of the subroutine STRIP_ZEROS. This
06EE 2437 : entry point has an additional longword, the return PC, on the stack on top
06EE 2438 : of the saved register array.
06EE 2439
06EE 2440 : 00(R0) - Return PC in mainline VAXSMULP or VAXSDIVP routine
06EE 2441 : 04(R0) - Saved R0
06EE 2442 : 08(R0) - Saved R1
06EE 2443 : etc.
06EE 2444 :-
06EE 2445
06EE 2446 MULP_BSBW_0:
06EE 2447 DIVP_BSBW_0:
50 04 C0 06EE 2448 ADDL #4,R0 ; Skip over return PC and drop into ...
06F1 2449
06F1 2450 :+
06F1 2451 : DIVP_0
06F1 2452 : MULP_DIVP_0
06F1 2453
06F1 2454 : There was nothing allocated on the stack other than the saved register
06F1 2455 : array when the access violation occurred. We merely pass control to common
06F1 2456 : code to restore the registers.
06F1 2457
06F1 2458 : 00(R0) - Saved R0
06F1 2459 : 04(R0) - Saved R1
06F1 2460 : etc.
06F1 2461 :-
06F1 2462
06F1 2463 DIVP_0:
06F1 2464 MULP_DIVP_0:
F90C' 31 06F1 2465 BRW VAX$DECIMAL_ACCVIO ; Join common code to restore registers
06F4 2466
06F4 2467 :+
06F4 2468 : DIVP_R6_R7
06F4 2469
06F4 2470 : An access violation occurred while one of the two input strings was being
06F4 2471 : converted to an array of longwords on the stack. The state of the stack
06F4 2472 : is rather complicated but R6 and R7 contain enough information to allow
06F4 2473 : the rest of the stack contents to be ignored.
06F4 2474
06F4 2475 : R6 - Count of longwords in quotient array on stack
06F4 2476 : R7 - Address of quotient array on stack
06F4 2477
06F4 2478 : 00(R0) - First longword of quotient array
06F4 2479
06F4 2480
06F4 2481 : zz-4(R0) - Last longword of scratch storage
06F4 2482 : zz+0(R0) - Digit count of dividend string
06F4 2483 : zz+4(R0) - Address of dividend string
06F4 2484 : zz+8(R0) - Saved R0
06F4 2485 : zz+12(R0) - Saved R1
06F4 2486 : etc.
06F4 2487
06F4 2488 : where zz = 4 * R6
06F4 2489 :-

```

50 08 A746 DE 06F4 2490
F904' 31 06F4 2491 DIVP_R6_R7:
06F4 2492 MOVAL 8(R7)[R6],R0 ; Discard everything on stack
06F9 2493 BRW VAX\$DECIMAL_ACCVIO ; Join common code to restore registers
06FC 2494
06FC 2495 END_MARK_POINT
06FC 2496
06FC 2497 .END

J 11
- VAX-11 Packed Decimal Arithmetic Instr 8-JAN-1985 17:27:01 VAX/VMS Macro V04-00
5-SEP-1984 00:44:34 [EMULAT.BUGSRC]VAXARITH.MAR;1 Page 58 (23)

...PC...	= 00000553		VAX\$ADD_PACKED_BYTE_R6_R7	00000170 RG 02
..ROPRAND..	= 000004BF	R 02	VAX\$DECIMAL_ACCVIO	***** X 00
ADDP4_B_DELTA_PC	= 00000003		VAX\$DECIMAL_EXIT	***** X 00
ADDP6_B_DELTA_PC	= 00000003		VAX\$DIVP	000004B1 RG 02
ADD_PACRED	= 000000E5	R 02	VAX\$MULP	000002AD RG 02
ADD_PACKED_BYTE_R6_R7	= 00000170	R 02	VAX\$REFLECT_TRAP	***** X 00
ADD_PACKED_BYTE_STRING	= 0000016A	R 02	VAX\$ROPRAND	***** X 00
ADD_SUBTRACT_EXIT	= 0000013D	R 02	VAX\$SUBP4	00000022 RG 02
ADD_SUB_24	= 000006BF	R 02	VAX\$SUBP6	00000000 RG 02
ADD_SUB_BSBW_0	= 000006C8	R 02		
ADD_SUB_BSBW_24	= 000006BC	R 02		
ADD_SUB_BSBW_4	= 000006C5	R 02		
ADD_SUB_V_ZERO_R4	= 0000001F			
ARITH_ACCVIO	= 0000068F	R 02		
CHECK_WRITE_ACCESS	= 00000292	R 02		
DECIMAL\$BINARY_TO_PACKED_TABLE	*****	X 00		
DECIMAL\$BOUNDS_CHECK	*****	X 00		
DECIMAL\$PACKED_TO_BINARY_TABLE	*****	X 00		
DECIMAL\$STRIP_ZEROS_R0_RT	*****	X 00		
DECIMAL\$STRIP_ZEROS_R2_R3	*****	X 00		
DECIMAL_ROPRAND	= 00000686	R 02		
DIVIDE_BY_ZERO	= 0000049E	R 02		
DIVP_0	= 000006F1	R 02		
DIVP_BSBW_0	= 000006EE	R 02		
DIVP_B_DELTA_PC	= 00000003			
DIVP_R6_R7	= 000006F4	R 02		
EXTEND_STRING_MULTIPLY	= 00000470	R 02		
HANDLER_TABLE_BASE	= 00000000	R 04		
MODULE_BASE	= 00000000	R 02		
MODULE_END	= 000006FC	R 02		
MULP_AT_SP	= 000006D4	R 02		
MULP_BSBW_0	= 000006EE	R 02		
MULP_B_DELTA_PC	= 00000003			
MULP_DIVP_0	= 000006F1	R 02		
MULP_DIVP_8	= 000006E8	R 02		
MULP_DIVP_R9	= 000006E0	R 02		
MULP_R8	= 000006CE	R 02		
MULTIPLY_DIVIDE_EXIT	= 0000037D	R 02		
MULTIPLY_STRING	= 0000066F	R 02		
PC_TABLE_BASE	= 00000000	R 03		
PS\$MN	= 00000008			
PSLSM_V	= 00000002			
PSLSM_Z	= 00000004			
PSLSV_CM	= 0000001F			
PSLSV_V	= 00000001			
PSLSV_Z	= 00000002			
QUOTIENT_DIGIT	= 000005CA	R 02		
SRMK_FLT_DIV_T	= 00000004			
STORE_RESULT	= 00000254	R 02		
SUBP4_B_DELTA_PC	= 00000003			
SUBP6_B_DELTA_PC	= 00000003			
SUBTRACT_PACKED	= 00000198	R 02		
SUB_PACKED_BYTE_R6_R7	= 0000022E	R 02		
SUB_PACKED_BYTE_STRING	= 00000228	R 02		
TABLE_SIZE	= 0000002A			
VAX\$ADDP4	= 00000030	RG 02		
VAX\$ADDP6	= 00000009	RG 02		

+-----+
! Psect synopsis !
+-----+

PSECT name

PSECT name	Allocation	PSECT No.	Attributes
. ABS .	00000000 (0.)	00 (0.)	NOPIC USR CON ABS LCL NOSHR NOEXE NORD NOWRT NOVEC BYTE
\$ABSS	00000000 (0.)	01 (1.)	NOPIC USR CON ABS LCL NOSHR EXE RD WRT NOVEC BYTE
VAX\$CODE	000006FC (1788.)	02 (2.)	PIC USR CON REL LCL SHR EXE RD NOWRT NOVEC LONG
PC_TABLE	00000054 (84.)	03 (3.)	PIC USR CON REL LCL SHR NOEXE RD NOWRT NOVEC BYTE
HANDLER_TABLE	00000054 (84.)	04 (4.)	PIC USR CON REL LCL SHR NOEXE RD NOWRT NOVEC BYTE

+-----+
! Performance indicators !
+-----+

Phase

Phase	Page faults	CPU Time	Elapsed Time
Initialization	75	00:00:00.20	00:00:02.24
Command processing	79	00:00:00.46	00:00:03.73
Pass 1	241	00:00:09.83	00:00:40.27
Symbol table sort	0	00:00:00.62	00:00:02.26
Pass 2	393	00:00:06.36	00:00:32.16
Symbol table output	0	00:00:00.08	00:00:00.54
Psect synopsis output	0	00:00:00.04	00:00:00.36
Cross-reference output	0	00:00:00.00	00:00:00.00
Assembler run totals	788	00:00:17.39	00:01:21.56

The working set limit was 1800 pages.

52274 bytes (103 pages) of virtual memory were used to buffer the intermediate code.

There were 20 pages of symbol table space allocated to hold 184 non-local and 116 local symbols.

2646 source lines were read in Pass 1, producing 25 object records in Pass 2.

23 pages of virtual memory were used to define 21 macros.

+-----+
! Macro library statistics !
+-----+

Macro Library name

\$255\$DUA18:[EMULAT.OBJ]VAXMACROS.MLB;1
 -\$255\$DUA18:[SYSLIB]STARLET.MLB;3
 TOTALS (all libraries)

Macros defined

12
6
18

318 GETS were required to define 18 macros.

There were no errors, warnings or information messages.

MACRO/LIS=LISS:VAXARITH/OBJ=OBJ\$:VAXARITH MSRC\$:VAXARITH/UPDATE=(BUGS:VAXARITH)+LIB\$:VAXMACROS/LIB

0441 AH-EF71A-SE
VAX/VMS V4.1 SRC LST MCRF UPD

