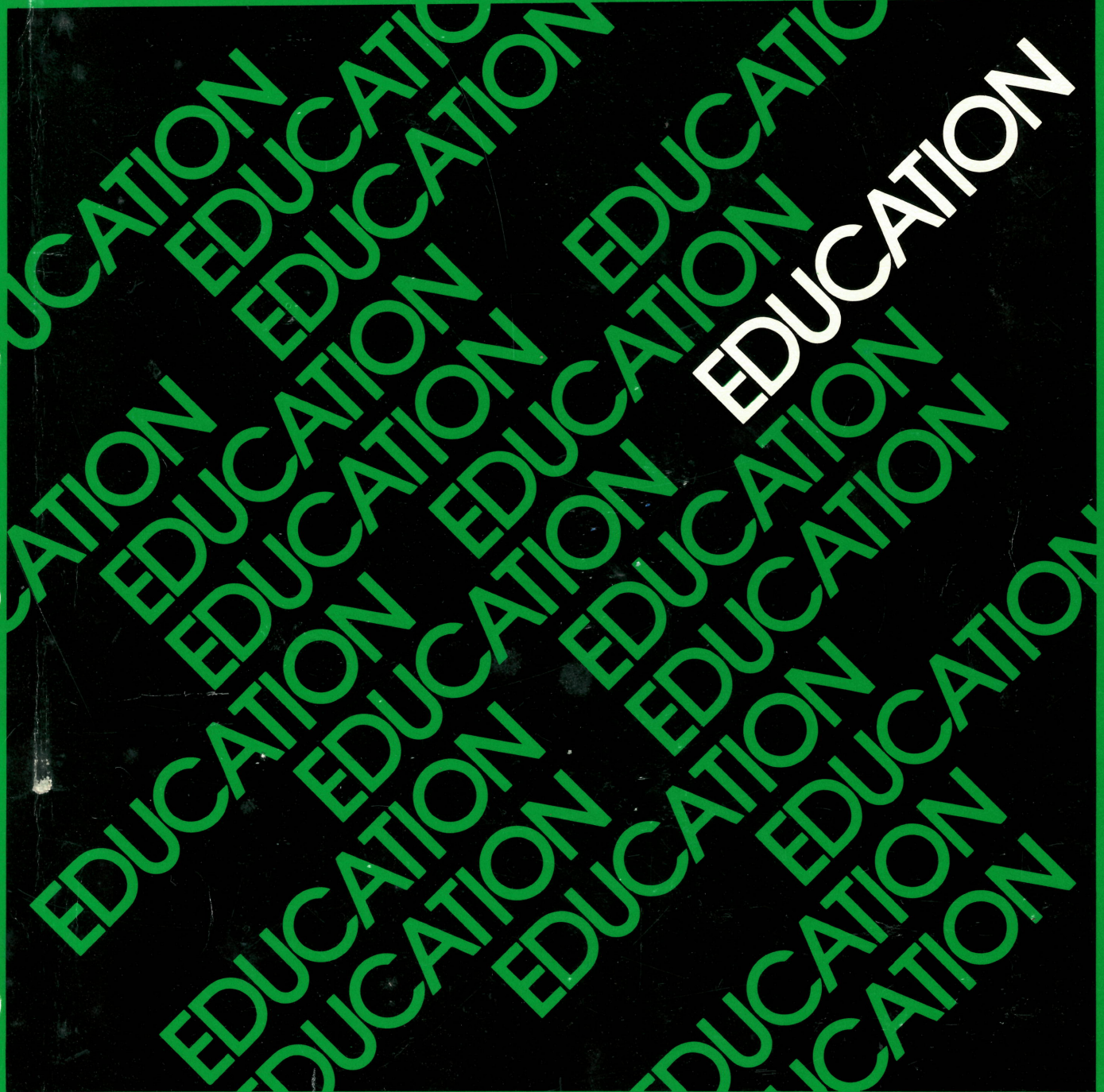


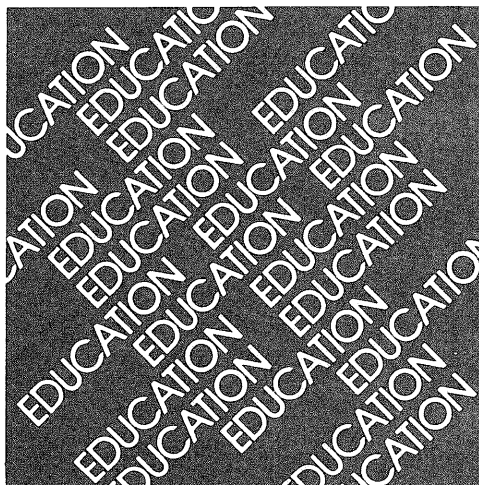
H. Bayhurst

digital

PDP-15
SYSTEM SOFTWARE
HANDOUTS



PDP-15
SYSTEM SOFTWARE
HANDOUTS



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PDP-15 SYSTEM SOFTWARE

COURSE ABSTRACT

This course is intended for programmers who wish to acquire a working familiarity with PDP-15 Assembly Language programming and the Disk or Advanced Monitor Operating Systems and the services provided by their monitors and associated system software. A portion of course time is devoted to supervised laboratory sessions.

PREREQUISITES

A working knowledge of the material presented in the *Introduction to Minicomputers* course.

COURSE OBJECTIVES

Upon successful completion of this course, the student will be able to:

- Write, run and modify assembly language programs using the PDP-15 instruction set and the MACRO assembler syntax.
- Interface programs to the Advanced or DOS I/O Monitor.
- Interact with the system by means of keyboard commands, system program command strings and programmed monitor requests.

COURSE OUTLINE

- I. Memory Organization; Memory and Addressing Modes
- II. PDP-15 Instruction Set
 - A. Memory Reference Instructions
 - B. Augmented Instruction Set

III. MACRO Assembler Syntax

IV. Tape and File Formats

V. Interrupt Systems

- A. Program Interrupt Control (PI)
- B. Automatic Priority Interrupt (API)

VI. System Programs

- A. EDIT, MACRO, LINKING LOADER, DDT
- B. PIP, PATCH, SGEN, UPDATE, CHAIN and EXECUTE

VII. I/O

- A. I/O Monitor
- B. System Macros

VIII. FORTRAN and MACRO Interface

IX. Handler Format

- X. Interaction with the system via keyboard commands and programmed monitor requests.

COURSE LENGTH

10 days

COURSE OUTLINE

WEEK I

MONDAY A.M.

- A. Description of Course
- B. Block Diagram of PDP-15
- C. Software Overview
- D. Memory Organization
 - 1. Addressing on the PDP-15
 - 2. Arithmetic on the PDP-15
- E. Central Processor Organization
- F. Introduction to the Instruction Set

MONDAY P.M.

- G. Memory Reference Instructions
- H. Operate Instructions
- I. EAE Instructions
- J. Introduction to MACRO-15

TUESDAY A.M.

- A. Review
- B. Subroutines
- C. Sum Group of Examples
- D. Indexed Instructions
- E. Absolute vs. Relocatable Programs

TUESDAY P.M.

- F. I/O Overview
- G. IOT Instructions and Dedicated I/O
- H. Paper Tape Formats

WEDNESDAY A.M.

- A. Console Description and Operation
- B. Operation of Disk Operating System
- C. Use of the Editor

WEDNESDAY P.M.

- D. LAB

THURSDAY A.M.

- A. MACRO-15 (again)
- B. Program Interrupt Facility (PI)
- C. Automatic Priority Interrupt (API)

THURSDAY P.M.

- D. LAB
- E. Program and Homework Review
- F. Quiz 1

FRIDAY A.M.

- A. Quiz 1 Review
- B. Overview of DOS Monitor
- C. File Formats and Directory Structure
- D. Overview of Monitor Supervised I/O

FRIDAY P.M.

- E. LAB

WEEK II

MONDAY A.M.

- A. Programmed I/O Commands
 - 1. Basic Operation
 - 2. System Macros
 - 3. DAT usage and Linking Loader
 - 4. User's Buffer Structure

MONDAY P.M.

- B. LAB

TUESDAY A.M.

- A. Linking Loader
- B. Libraries and UPDATE
- C. CHAIN and EXECUTE
- D. PIP

TUESDAY P.M.

- E. LAB

WEDNESDAY A.M.

- A. Fortran and Macro Interface
- B. DDT
- C. ↑Q
- D. DUMP

WEDNESDAY P.M.

- E. LAB
- F. Program and Homework Review
- G. Quiz 2

THURSDAY A.M.

- A. I/O Handler Format
- B. Sample Handler

THURSDAY P.M.

- C. LAB
- D. Final Exam

FRIDAY A.M.

- A. Monitor and System Modification
 - 1. PATCH
 - 2. SGEN
- B. Unichannel Discussion
- C. Final Exam Review

FRIDAY P.M.

- D. Optional Lab

BLOCK DIAGRAM OF THE PDP-15 SYSTEM

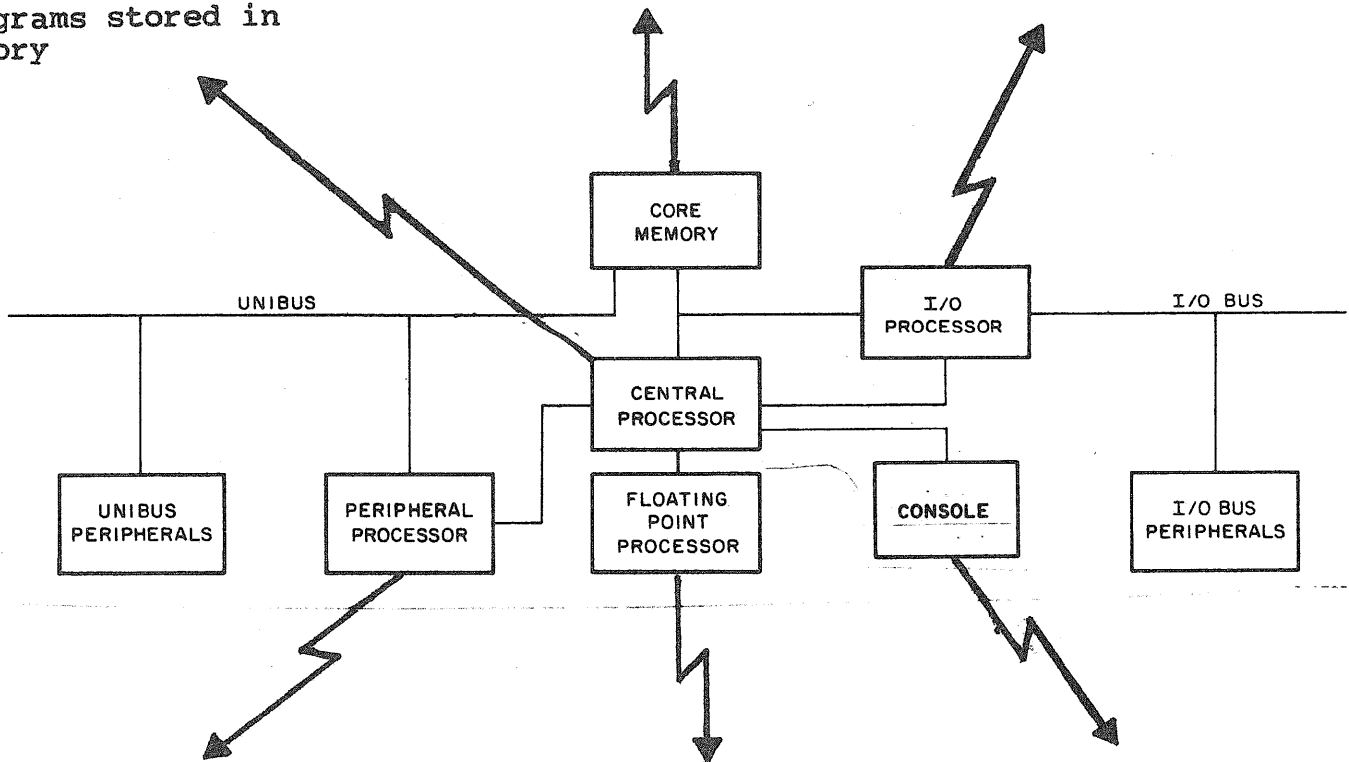
Computer: a machine which inputs data from the outside world, processes the data, perhaps puts it in temporary storage and finally outputs the data and/or results in the form of hard copy, displays or commands to other devices it may be controlling.

Program: a sequence of instructions to a computer, specifying the necessary steps to solve a problem (sometimes it includes the data it is to work on).

- main component
- handles bidirectional communication with memory and I/O processor
- performs arithmetic and logical operations
- controls and executes programs stored in memory

- main storage area within computer from which instructions are fetched and executed (data may also be stored here)
- various types of memory available

- handles peripheral data transfers
- coordinates transfer between: CP & peripherals memory & peripherals



- second general purpose processor
- second I/O bus allowing use of PDP-11 peripherals

- allows use of floating point arithmetic via over 100 instructions without use of complex software routines

- allows operator communication in system:
 - starting & halting of programs
 - monitoring of registers
 - modification of memory

SOFTWARE

To utilize the powerful PDP-15 hardware a number of operating systems have been developed (special applications packages are also available). See chapter 10 of the System Reference Manual for descriptions of each of these systems.

DOS-15, DISK OPERATING SYSTEM

Disk Operating System (DOS-15) is an integrated set of software designed to meet the demands of research, engineering, and industrial environments. It includes the software necessary for simplified programming and efficient operations. DOS-15 brings to the user the advantage of disk resident storage via rapid access to the system's resources.

The DOS Monitor, the heart of the system, incorporates all the functions of the "Advanced Software System" plus the added power of fully automatic random access file operation. The user controls the operating system by instructions to the Monitor. The Monitor runs the jobs, supervises data and file manipulation, and interacts with the operator/user in a simple conversational manner.

Noteworthy features of DOS-15 are:

Disk Resident System Software

All DOS-15 System Software resides on either DECdisk, or RP15 Disk Pack, or RK15 disk cartridge.

Interactive Operation

An interactive keyboard/program Monitor permits device-independent programming, and automatic calling and loading of system and user programs.

Conversational Mode

System Utility Programs interact with the operator/user in a simple, conversational manner.

Programmed Monitor Commands

Input/Output programming is simplified by the use of a set of system commands which are standardized for system-supported I/O devices.

I/O Device Handlers

Data and file manipulating I/O device handlers are supplied for standard system peripherals, allowing device independence and overlapped computation, and I/O.

User-Created System Files

The user may easily incorporate his own software into the operating system, thereby tailoring the system to his hardware and software needs.

Programming Languages

FORTRAN IV, FOCAL, and MACRO-15 programming languages are offered.

Bank and Page Modes

Choice of 8K (Bank Mode) or 4K (Page Mode) direct addressability. Page Mode operation permits modification via the index register.

Disk File Structure

The disk file structure allows the most efficient use of disk capacity and data retrieval for processing via:

System supported DECdisk, Disk Packs, and Disk Cartridge Devices, providing both economy and storage capacity.

Virtually unlimited data capacity (Disk Pack = 83.7 million words, DECdisk = 2.09 million words, Disk Cartridge = 9.6 million words). Random/Sequential File Access furnishes file protection through unique user directories and associated user identification codes. Files can be made invisible to other users, but with privileged access via a supervisory code.

User/user file independence—identically named unformatted Input/Output (FORTRAN IV).

Random Access-formatted as well as unformatted Input/Output (FORTRAN IV).

Dynamic Storage Allocation

The available disk storage is automatically allocated for optimum storage utilization.

Dynamic Buffer Allocation

Input/Output core is automatically optimized by the Monitor. It allocates only that space which is required for the system and the user.

Batching Operation

An alternative to interactive operation is a batching mode which permits the sequencing of console commands to come from paper tape or cards.

Input/Output Spooling

DOS-15 systems using the RK15/RK05 Unichannel Disk System, provides spooling of card reader, line printer, and XY plotter data.

Spooling is a method of storing (queueing) data to and from slow speed devices on the high speed RK05 disk. This dramatically improves system performance.

Spooling is only provided for devices interfaced to the UNIBUS of the RK15 Disk System (i.e., the CR11, LP11, LS11, and XY11). Spooling requires 8K of local PDP-11 memory.

The following software is available as part of DOS-15:

Monitors

Resident Monitor
Keyboard Command Decoder
Batch Processor
System Loader
PIREX (Peripheral Processor RK15 Only)

Languages

FORTRAN IV (F4X, FPPF4X)
FOCAL
MACRO-11 (Assembler RK15 Only)
MACRO-15 (Assembler)
ALGOL (optional)

Text Editors

EDIT
EDITVP (Storage Scope Editor)
EDITVT (Graphic Display Editor)

Loaders

Linking Loader
CHAIN & EXECUTE (Overlay Loaders)
ABS 11 (RK15 Only)

Debuggers

DDT (Dynamic Debugging Technique)
DUMP (Core Dump Lister)
QFILE (Store/Retrieve Core Dumps)

Utilities (General)

DTCOPY (DECtape Copier)
MTDUMP (Magtape Utility)
PIP (Peripheral Interchange Program)
SRCCOM (Source Compare)
UPDATE (Library File Manager)
8TRAN (PDP-8 to PDP-15 Translator)
89TRAN (PDP-8 to PDP-9 Translator)
TKB (RSX-15 Task Builder)

Utilities (System)

DOSSAV (Disk Save/Restore)
RFBOOT (DECdisk Bootstrap)
RPBOOT (Disk Pack Bootstrap)
RKBOOT (Disk Cartridge Bootstrap)

I/O Handlers

CDB (Card Reader for CR03B, CR15 or CR11)
DOSBCD (Batch Card Reader)
DKA, DKB, DKC, DKL, (RF15/RS09 DECdisk)
DPA, DPB, DPC, DPL (RP15/RP02 Disk Pack)
RKA, RKB, RKC, RKL (RK15/RK05 Disk Cartridge)
DTA, DTC, DTD, DTE, DTF, (DECtape)
LKA (LK35 Graphics Keyboard)
LPA (Line Printer for LP15, LS11 or LP11)
LVA (Line Printer/Plotter)
MTA, MTC, MTF (Magtape)
PPA, PPB, PPC (Paper Tape Punch)
PRA, PRB (Paper Tape Reader)
TTA (Teletype)
VPA (Storage Scope)
VTA (VT15 Graphic Display)
VWA (VW01 Writing Tablet)
XYA (XY11 Plotter, RK DOS only)

Checkout-Package

RF.CHK (DECdisk Checkout)
RP.CHK (Disk Pack Checkout)
RK.CHK (Disk Cartridge Checkout)

Minimum Hardware

KP15 Central Processor
16,384 18-bit Core Memory
Console Terminal
PC15 High Speed Paper Tape Reader and Punch
KE15 Extended Arithmetic Element
TC15 DECtape Control¹—or TC59 Magtape Control
1 TU56 Dual DECtape Transport—or 1 TU10, TU20,
or TU30 (7 or 9 track) Magtape Transport
RK15 DECdisk Control or RP15 Disk Pack Control or
RK15 System
1 RS09 Disk Drive or 1 RP02 Disk Pack Drive or
1 RK05 Drive

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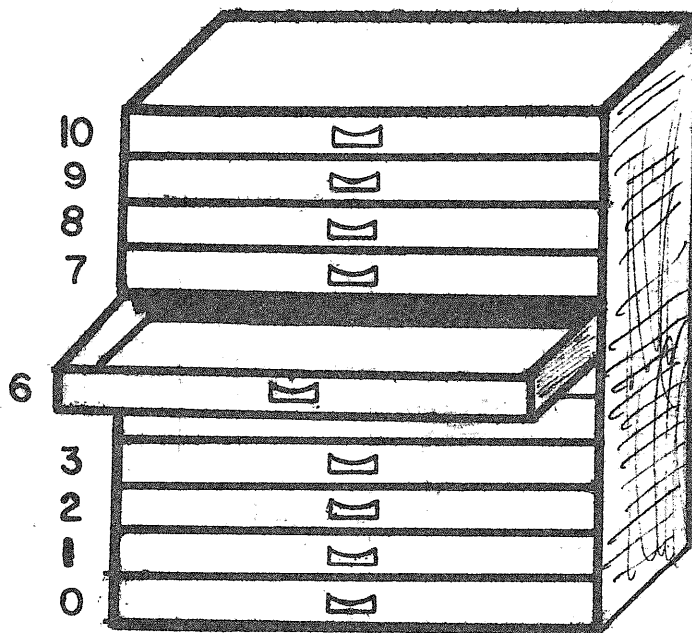
MEMORY

Memory: the main storage area for computer instructions and system data.

In order for a program to be executed, it must be placed ("loaded") into memory.

Memory is storage space--a place to keep things for a while. It can hold either data or instructions.

Memory, often referred to as main storage, is much like a large chest of drawers.



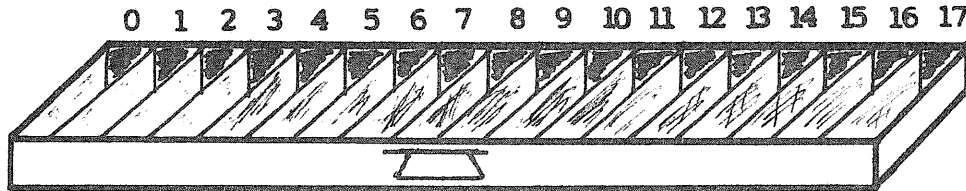
MEMORY

- a. You can store something in each drawer.
- b. You must examine a basic storage unit when looking for something -- in the chest this unit is a drawer; in a computer this unit is a location (word).
- c. You refer to these basic units by numbers - you tell someone to look for something in the 3rd drawer from the bottom; you tell the processor to look for something in location 3.

In a computer, the numbers of the locations are called addresses. Address numbers begin with zero as shown above.

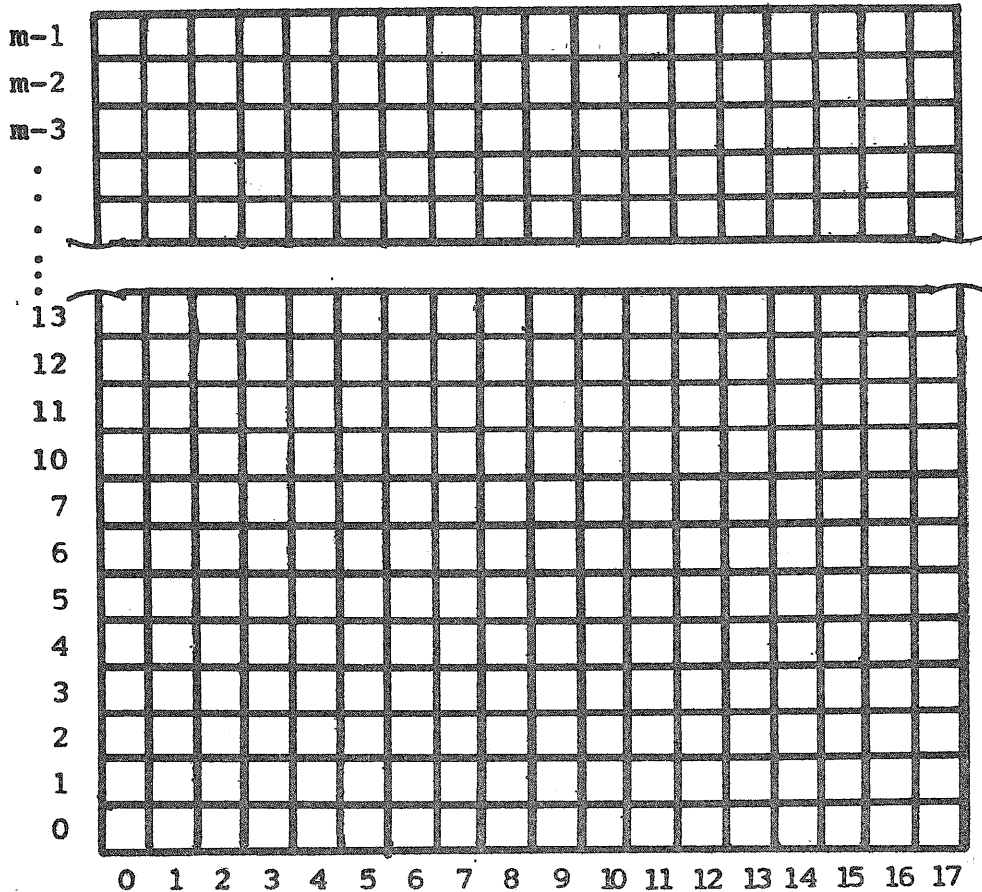
Each location or word in a PDP-15 is partitioned into smaller subdivisions called bits. Bits are subdivisions which have binary values, either 1 or 0.

Locations subdivided into bits are like drawers partitioned into small slots. To examine a bit, you must first look at the entire location; to look into a slot you must first pull out the entire drawer.



LOCATION

The bits in a location are also numbered beginning with zero, but these numbers are not generally referred to as addresses. So you may imagine memory as a chest of numbered drawers, each containing numbered slots; or you may simplify your model and imagine a matrix of m numbered locations, each containing 18 numbered bits as shown below.



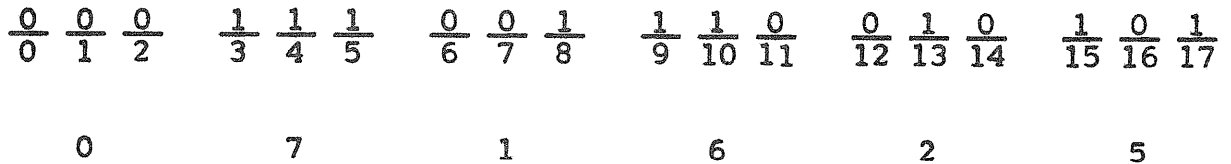
WORD LENGTH

The PDP-15 has an 18 bit word.

The bits are labelled from left to right, starting with 0 and ending with 17:



Because one octal digit is the equivalent of three binary digits, the contents of an 18 bit word is often given as a string of 6 octal digits:



Binary value: 000111001110010101

Octal value: 071625

SIGNED NUMBERS

Words may be looked at as instructions or as data (numerical values).

When viewed as numerical values, numbers are usually looked at as being SIGNED.

The sign of a number is determined by the most significant bit:

- MSB=0, the number is POSITIVE.
- If
- MSB=1, the number is NEGATIVE.

EXAMPLES

Positive numbers --

010 111 001 011 101 000	OR	271350
000 001 101 110 010 100	OR	015624
000 000 000 000 000 001	OR	000001
000 000 000 000 000 000	OR	000000

Note that "0" (zero) is considered a positive number because bit 0, its most significant bit, is 0.

Negative numbers --

101 111 001 100 010 001	OR	571421
111 000 000 010 110 011	OR	700263
111 111 111 111 111 111	OR	777777
100 000 000 000 000 000	OR	400000

Range --

<u>POSITIVE</u>	<u>NEGATIVE</u>
000000	
000001	777777
000002	777776
⋮	⋮
100000	700000
⋮	⋮
200000	600000
⋮	⋮
300000	500000
⋮	⋮
377776	400002
377777	400001
	400000

COMPLEMENT NUMBERS

All negative numbers on the PDP-15 are expressed in COMPLEMENT FORM rather than sign-magnitude form.

sign magnitude

+5	000 000 000 000 000 101	OR	000005
-5	100 000 000 000 000 101	OR	400005

There are two forms of complement numbers used on the PDP-15:

1's COMPLEMENT -- the 1's complement of a binary number is the result of inverting each bit position.

ex.

number	010 101 001 000 100 011	OR	251043
1's complement	101 010 110 111 011 100	OR	526734

Note that when you add a number to its 1's complement, the sum is:

111 111 111 111 111 111	OR	777777.
-------------------------	----	---------

Hence, an easy way to compute the 1's complement of an octal number is to subtract it from 777777.

777777
- 251043
526734

2's COMPLEMENT -- the 2's complement of a binary number is the result of adding "1" to the 1's complement.

i.e. (the 2's complement) = (the 1's complement) + 1.

ex.

number	010 101 001 000 100 011	OR	251043
1's complement	101 010 110 111 011 100	OR	526734
2's complement	101 010 110 111 011 101	OR	526735

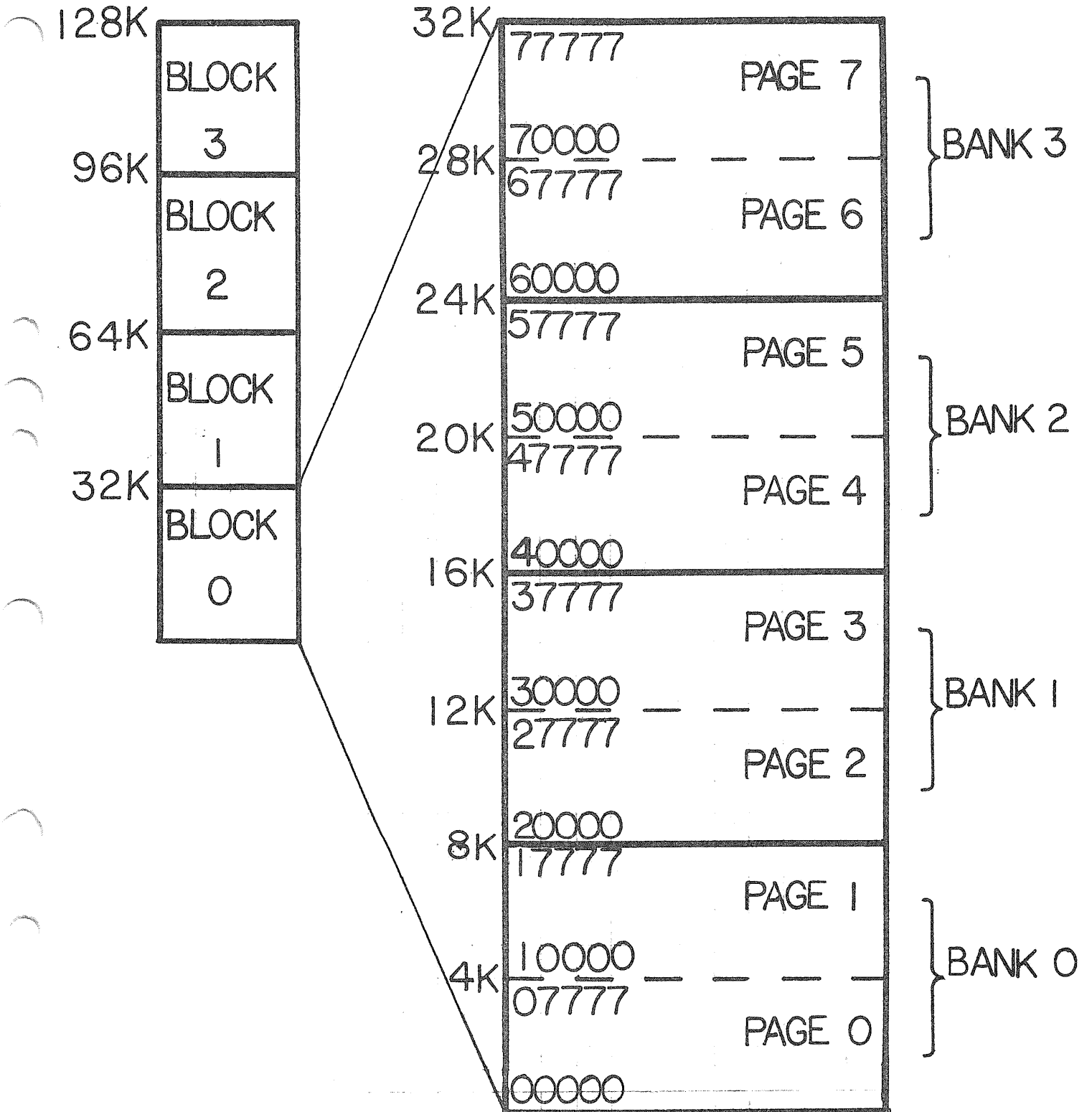
Note that an easy way to compute the 2's complement of an octal number is to subtract it from 777778 or a similar value value where the 8 is in the rightmost non-zero position of the number being complemented.

251043	777778	BUT	251040	777780
	-251043			-251040
	526735			526740

Note that when you add a number to its 2's complement, the sum is: 000000.

251043	251040
+526735	+526740
000000	000000

MEMORY ORGANIZATION



NOTE: There are 10000 (octal) locations in each page.
 Each page starts with an address that is a multiple of 10000.

There are 20000 (octal) locations in each bank.
 Each bank starts with an address that is an even multiple of 20000.

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DOUBLE PRECISION ADDITION

/ Program to illustrate the use of two's complement addition in performing
 / a double precision add. Two words are used for each DP number. The
 / first word of the pair contains the sign and the most significant part.
 / The second word contains the low order part. It is important to note
 / that all bits of the low order part are numeric bits, i.e., the sign
 / bit is considered as a numeric bit, not as a sign. For illustration
 / purposes, assume that the word size is only six bits. Some DP numbers
 / follow:

<u>NUMBER</u>	<u>HIGH</u>	<u>LOW</u>
0025	000 000	010 101
0063	000 000	110 011
0377	000 011	111 111
-1	111 111	111 111
-100	111 111	000 000
-40	111 111	100 000

The program to perform DP addition follows:

```

CLA!CLL      /Clear AC and link
TAD AL       /Get low half of A
TAD BL       /Add low half of B
DAC CL       /Save low half of result
CLA          /Clear AC, but link remains
GLK          /If addition of AL and BH caused a carry, link=1
              /and the one is placed in AC (17) so that it can
              /be added to the high half. If no carry, AC 17=0
TAD AH       /Add in high order parts of the two numbers
TAD BH
DAC CH
  
```

/ EXAMPLE 1. DP ADD 127+ 306 = 735

	<u>HIGH</u>	<u>LOW</u>	
/ 127=	000 001	010 111	
/ 306=	000 011	000 110	
/ Add low order		011 101	Link = 0
/ Add high order	000 100		
/ Add in LINK	0		
/ Result	000 100	100 101	= 435

/ EXAMPLE 2. DP ADD 1254 + 2362 = 3636

	<u>HIGH</u>	<u>LOW</u>	
/ 1254=	001 010	101 100	
/ 2362=	010 011	110 010	
/ Add low order		011 110	Link = 1
/ Add high order	011 101		
/ Add in LINK	1		
RESULT	011 110	011 110	= 3636

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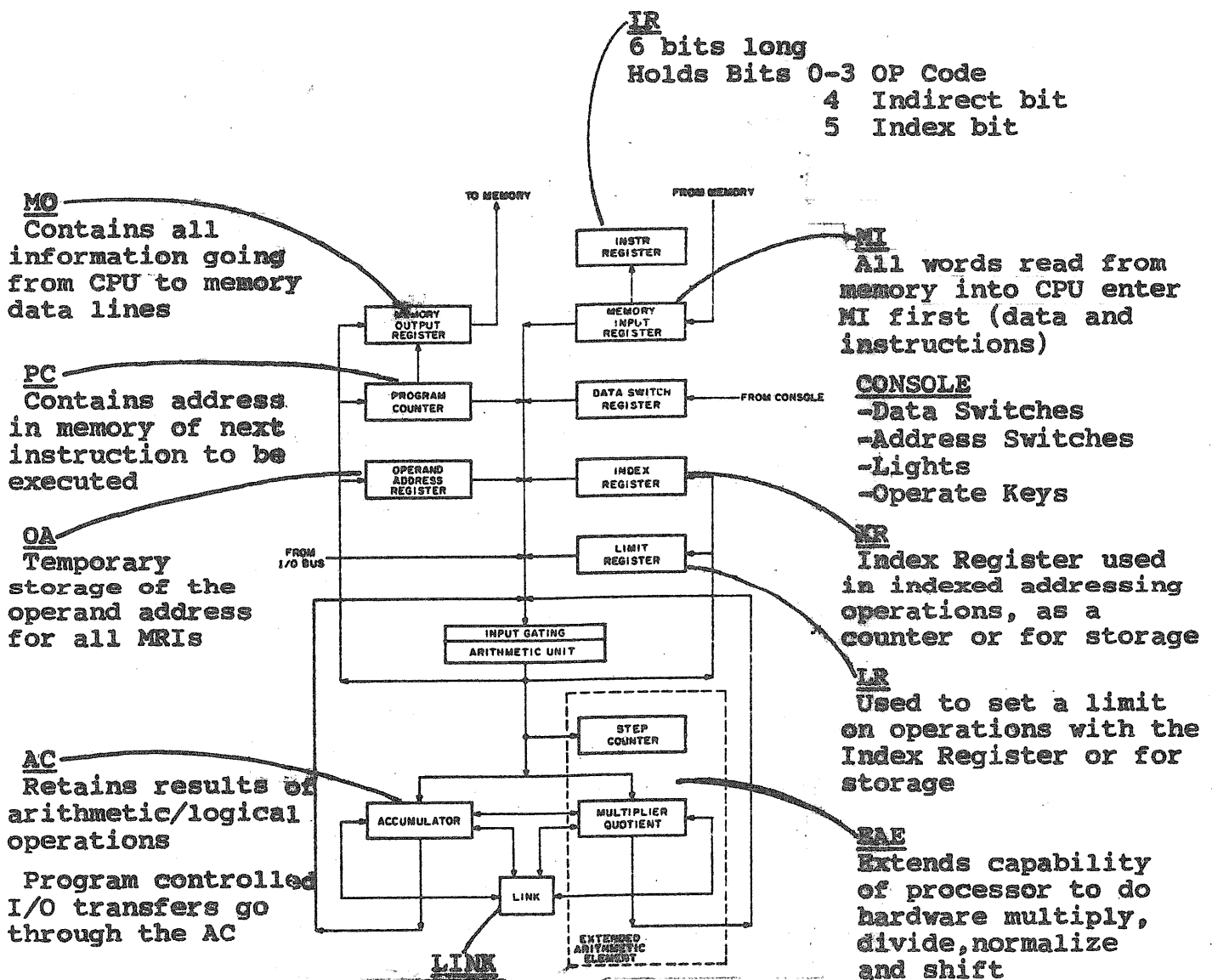
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CPU

CENTRAL PROCESSOR: the Central Processor Unit (CPU) functions as the main component of the computer by carrying on bidirectional communication with both the memory and I/O Processor. Provided with the capability to perform all required arithmetic and logical operations, the central processor controls and executes stored programs. It accomplishes this with an extensive complement of registers, control lines, and logic.



MO
Contains all information going from CPU to memory data lines

PC
Contains address in memory of next instruction to be executed

OA
Temporary storage of the operand address for all MRIs

AC
Retains results of arithmetic/logical operations

Program controlled I/O transfers go through the AC

IR
6 bits long
Holds Bits 0-3 OP Code
4 Indirect bit
5 Index bit

MI
All words read from memory into CPU enter MI first (data and instructions)

CONSOLE
-Data Switches
-Address Switches
-Lights
-Operate Keys

IR
Index Register used in indexed addressing operations, as a counter or for storage

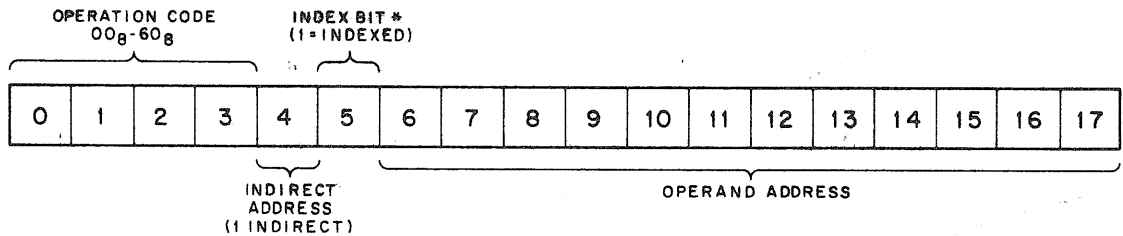
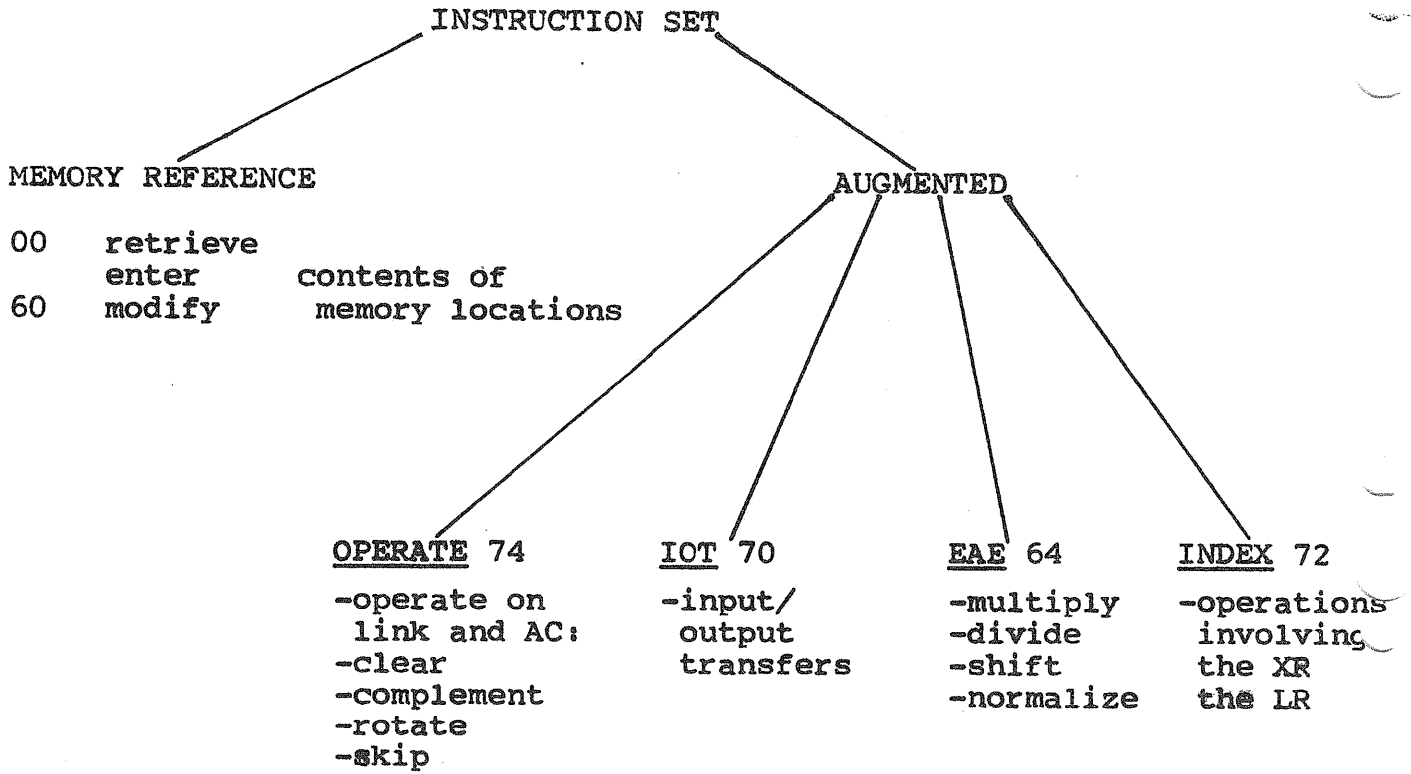
IR
Used to set a limit on operations with the Index Register or for storage

BAE
Extends capability of processor to do hardware multiply, divide, normalize and shift

LINK
1 bit extension of the AC

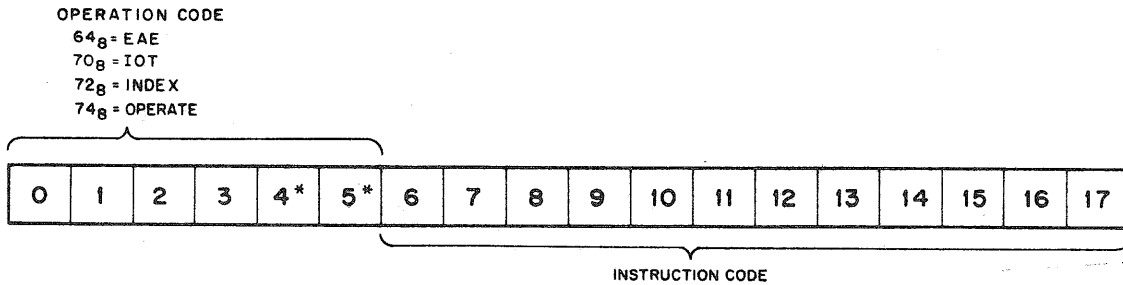
CPU Options

- Memory Protect
- Memory Protect & Relocate
- Powerfail



*USED AS A THIRTEENTH ADDRESS BIT IN BANK MODE

Memory Reference Instruction Word



*THESE BITS USED AS PART OF THE INSTRUCTION CODE IN EAE AND OPERATE INSTRUCTIONS

Augmented Instruction Format

.TITLE MACRO ASSEMBLY LANGUAGE EXAMPLE

00100

.ABS

.LOC 100

00100 600100
00101 600100
00102 000100

START JMP .
JMP START
START

00103 140103
00104 200103
00105 140000
00106 000103
00107 000103

DZM DZM DZM
LAC DZM
DZM
DZM
.DSA DZM

/VARIABLES

00110 201025
00111 041025
00112 201026
00113 041026

LAC CNT#
DAC CNT#
LAC CNT2#
DAC CNT2

/LITERALS

00114 201031
00115 201024
00116 201032
00117 201033
00120 201034
00121 001034
000200
00122 201036

A LAC (1
LAC TABLE
LAC (TABLE
LAC (A
LAC (C
(C
B=200
LAC (B

/INDIRECT

00123 200100
00124 220100
00125 220100
00126 021005
00127 600100
00130 600100
00131 021036

LAC 100
LAC+ 100
LAC+ START+2
JMP+ A1
JMP 100
JMP START
JMP+ (100

.EJECT


```

/
/
01005          .LOC 1005
01005 000114   A1      A
          000100   B=100
          000004   C=4
M 01006 000100   D      B
01007 000004   C
01010 200100   LAC B
01011 200114   LAC A
01012 200220   LAC B+C+A
01013 201037   LAC (B+C+A)
01014 201040   LAC (E#)
/
DM 01015 001000   D      D      /MULTI-DEFINED SYMBOL
U 01016 001030   DD     /UNDEFINED SYMBOL
          000110   B=110
          01017 000110   B
          01020 200110   LAC B
          01021 201041   LAC (B)
/
G 01022 000010   A=B+C
01023 001042   A=(B+C)
/
01024 000000   TABLE 0
          000100   .END START
01031 000001   *L
01032 001024   *L
01033 000114   *L
01034 000004   *L
01035 000200   *L
01036 000100   *L
01037 000220   *L
01040 001027   *L
01041 000110   *L
01042 777074   *L
          SIZE=01044
          4 ERROR LINES

```

.TITLE PSEUDO OPS EXAMPLES

.ABS

00100

.LOC 100

00100 200100

LAC 100

/*
/.REPT

.REPT 4
RTL

00101 742010

00102 742010

00103 742010

00104 742010

*R
*R
*R

.REPT 5,1
DZM BUFF

00105 140112

00106 140113

00107 140114

00110 140115

00111 140116

ZERO

*R
*R
*R
*R

/*
/RESERVING STORAGE

00112

00117 000000

BUFF .BLOCK 5
BUFFE 0

00120 000000

00126 000000

00126 000000

TABLE 0
.LOC +5
TABLEE 0

/*
/CONDITIONALS

.IFDEF A
LAC A+1
.ENDC

000002

B=2

.IFPOZ B
TAD (B
.ENDC

00127 341123

/*
/PAGING THE LISTING

.EJECT

/*
/CHANGING THE LOCATION COUNTER
/*****L ****
/

01050

.LOC 1050

/*
/.ASCII AND .SIXBT
/*****
/

01050 405010
01051 251236
01052 456131
01053 606424
01054 202532
01055 020232
01056 426472
01057 340616
01060 425020
01061 000000

MESS .ASCII /A BROKEN/<15><12>

.ASCII / UP MESSAGE! /

01062 406050
01063 330544
01064 314000
01065 000000

MESS1 .ASCII 'ABC123'

01066 010203
01067 616263

NAME1 .SIXBT 'ABC123'

01070 406050
01071 330400

MESS2 .ASCII 'ABC1'

01072 010203
01073 610000

NAME2 .SIXBT 'ABC1'

01074 406050
01075 330544
01076 314000
01077 000000

MESS3 .ASCII 'AB'/C123/

01100 010203
01101 616263

NAME3 .SIXBT 'AB'/C123/

/*
/PAGE EJECT ALSO CAUSED BY .TITLE
/*****
/

.TITLE PSEUDO OP EXAMPLES--THIS PAGE: MACROS

/*
/DEFINING A MACRO

/*
/

.DEFIN SUB A,B,C
LAC B
TCA
TAD A
DAC C
.ENDM

/*
/CALLING A MACRO

/*
/

			SUB BUFPE,BUFF,TABLE
01102	200112	*G	LAC BUFF
01103	740031	*G	TCA
01104	340117	*G	TAD BUFPE
01105	040120	*G	DAC TABLE

			SUB BUFPE,(5,TABLE+1
01106	201124	*G	LAC (5
01107	740031	*G	TCA
01110	340117	*G	TAD BUFPE
01111	040121	*G	DAC TABLE+1

			SUB (MESS1,(MESS,TABLE+2
01112	201125	*G	LAC (MESS
01113	740031	*G	TCA
01114	341126	*G	TAD (MESS1
01115	040122	*G	DAC TABLE+2

			SUB BUFPE,TABLE,TEMP#
01116	200120	*G	LAC TABLE
01117	740031	*G	TCA
01120	340117	*G	TAD BUFPE
01121	041122	*G	DAC TEMP#

			.END
01123	000002	*L	
01124	000005	*L	
01125	001050	*L	
01126	001062	*L	

SIZE=01127

NO ERROR LINES

)

)

)

)

)

)

)

)

```

*****
*
*   SUM   GROUP   *
*
*****

```

THERE ARE 4 PROGRAMS IN THIS SERIES. THEY OFFER SOLUTIONS (EACH USING A DIFFERENT ADDRESSING MODE) TO THE SAME PROBLEM:

TABLES A, B AND C EACH CONTAIN 5-ONE WORD ENTRIES. ADD CORRESPONDING ENTRIES FROM TABLES A AND B AND STORE THE RESULT IN THE CORRESPONDING ENTRY IN TABLE C, I.E. $C(I) = A(I) + B(I)$. DO THIS WITHOUT CHANGING ANY OF THE VALUES IN TABLES A AND B.

THESE PROGRAMS WERE WRITTEN TO ILLUSTRATE THE VARIOUS TYPES OF ADDRESSING AVAILABLE ON THE PDP-15 AND TO DEMONSTRATE CERTAIN MACRO LANGUAGE ELEMENTS AND ASSEMBLER DIRECTIVES:

- (2-11;2-15) *STATEMENT FORMAT
- (3-3) *TYPE OF BINARY OUTPUT TO BE GENERATED
 - .ABS
 - .ABSP
- (3-10) *SETTING THE LOCATION COUNTER
 - .LOC
- (3-2) *GETTING HEADINGS ON ASSEMBLY LISTINGS
 - .TITLE
- (3-10) *SPECIFYING WHETHER NUMBERS ARE OCTAL OR DECIMAL
 - .DEC
 - .OCT
- (3-11) *RESERVING BLOCKS OF MEMORY FOR STORAGE
 - .BLOCK
- (2-6) *VARIABLES
 - USING #
- (2-13) *LITERALS
 - USING ()
- (2-3)
- (2-15) *DEFINING THE VALUE OF SYMBOLS
 - USING THE SYMBOL AS A LABEL
 - USING DIRECT ASSIGNMENTS
- (2-6)
- (2-16) *STORING VALUES IN SUCCESSIVE LOCATIONS
- (3-11) *SPECIFYING THE PHYSICAL END OF PROGRAM
 - .END

THE ABOVE IS JUST A SELECTED LIST OF ASSEMBLY LANGUAGE ELEMENTS AND ASSEMBLER DIRECTIVES (ALSO CALLED PSEUDO OPERATIONS). MAKE SURE YOU ARE FAMILIAR WITH THEM. EACH TOPIC IS DISCUSSED IN THE MACRO MANUAL STARTING ON THE PAGE GIVEN IN PARENTHESES.


```

.TITLE SOLUTION USES AN ADDRESS MODIFICATION TE
/
.ABS /ABSOLUTE PROGRAM; LOADED BY THE
/ABSOLUTE LOADER; TO RUN IN BANK MODE
00100 .LOC 100 /SET LOCATION COUNTER TO 100
/
00100 777773 START LAW -5 /SET UP NEGATIVE COUNTER OF -5
00101 040133 DAC COUNT /AS EACH TABLE HAS 5 ENTRIES.
/
00102 200113 AA LAC A /PERFORM ADDITION
00103 340120 BB TAD B /OF ENTRIES
/
00104 040125 CC DAC C /STORE RESULT IN C.
/
00105 440102 ISZ AA /MODIFY LOCATIONS AA, BB AND CC
00106 440103 ISZ BB /SO THAT THEY REFERENCE THE
00107 440104 ISZ CC /NEXT LOCATION IN THE TABLE.
/
00110 440133 ISZ COUNT# /ARE WE DONE?
/
00111 600102 JMP AA /NO! -- BECAUSE WE DIDN'T SKIP
/GO BACK FOR MORE.
/
00112 740040 HLT /YES! DONE.
/NOTE! IF THIS PROGRAM IS TO BE
/RUN AGAIN AA, BB AND CC SHOULD
/BE REINITIALIZED (CAN YOU THINK
/OF WAYS THIS COULD BE DONE?)

/SET UP TABLES
/
00113 000001 A 11 21 31 41 5
00114 000002
00115 000003
00116 000004
00117 000005
00120 000002 B 21 21 21 21 2
00121 000002
00122 000002
00123 000002
00124 000002

/JUST RESERVE SPACE FOR TABLE C
/
00125 C .BLOCK 5
/
00132 000000 ENDLOC 0
/
/THIS IS THE END OF THE PROGRAM.
/THEREFORE, THIS IS WHERE WE WANT TO PUT
/THE ".END" STATEMENT.
/
000100 .END START
SIZE=00134 NO ERROR LINES

```


.TITLE COMMENTARY ON SUM1

 * POINTS TO BE NOTED *

/*1. THIS PROGRAM USES INDIRECT ADDRESSING THROUGH
 /LOCATIONS AA, BB AND CC. LOCATION AA CONTAINS THE
 /ADDRESS OF TABLE A. NOTE THAT THIS IS ACCOMPLISHED
 /SIMPLY BY PLACING THE SYMBOL "A" IN THE OPERATOR
 /FIELD. WHEN THE ASSEMBLER EVALUATES THIS STATEMENT,
 /IT LOOKS THROUGH ITS SYMBOL TABLES FOR THE VALUE OF
 /SYMBOL "A". ***NOTE THAT THE VALUE OF THE SYMBOL "A"
 /IS 113, THE ADDRESS OF THE LOCATION IT NAMES----
 /NOT THE CONTENTS OF THE LOCATION IT NAMES. THIS IS
 /A VERY IMPORTANT DISTINCTION!!

/ SIMILARLY, LOCATION BB CONTAINS THE ADDRESS OF
 /TABLE B BECAUSE THE SYMBOL "B" IS GIVEN AS THE CONTENTS
 /OF LOCATION BB. THE VALUE OF THE SYMBOL B IS THE
 /ADDRESS OF THE LOCATION IT NAMES, 120, NOT THE CONTENT
 /OF THAT LOCATION. SIMILARLY FOR LOCATION CC. THE
 /CONTENT OF LOCATION CC IS GIVEN AS "C". THE VALUE OF
 /THE SYMBOL "C" IS 130, THE ADDRESS OF THE LOCATION
 /IT NAMES. HENCE, THE CONTENT OF LOCATION CC IS 130.

/*2. NOTE THAT DATA IS SEPARATED FROM THE INSTRUCTIONS
 /WHICH OPERATE ON THE DATA. WE DO NOT WANT TO "FALL INTO
 /THE DATA AND START EXECUTING IT. WHAT INSTRUCTION WOULD
 /WE HAVE IF WE "FELL INTO" AND TRIED TO EXECUTE LOC 113"

/*3. NOTE THE USE OF "COUNT" AS A VARIABLE, (WHAT IS IT
 /IN THE USE OF THE SYMBOL "COUNT" THAT MAKES IT A
 /A VARIABLE?). WHAT LOCATION IS SET UP BY THE ASSEMBLER
 /SO THAT IT MAY BE REFERENCED USING "COUNT"?

/*4. NOTE THAT WE RESERVE SPACE FOR TABLE C EVEN THOUGH
 /IT IS AT THE VERY END OF THE WRITTEN PROGRAM. WHAT
 /PROBLEM WOULD WE RUN INTO IF WE DID NOT RESERVE THIS
 /SPACE? (HINT: SEE NOTE *3)

/*5. THINK ABOUT THE "LAC* AA" INSTRUCTION IN LOC 102.
 /THIS SAYS: LOAD THE ACCUMULATOR WITH THE CONTENT OF
 /THE LOCATION POINTED TO BY LOCATION AA. THE FIRST TIME
 /THRU, LOCATION AA CONTAINS 113, THE ADDRESS OF THE FIRST
 /LOCATION IN TABLE A. IN LOCATION 105, THERE IS AN
 /"ISZ AA", I.E. INCREMENT THE CONTENT OF LOCATION AA BY 1
 / (SO THAT AA NOW POINTS TO THE NEXT ENTRY IN TABLE A).
 /THUS, THE NEXT TIME "LAC*AA" IS EXECUTED, THE AC IS
 /LOADED WITH THE NEXT ENTRY FROM TABLE A. THIS METHOD
 /IS VERY NICE FOR STEPPING THROUGH TABLES. NOTE, HOWEVER,
 /THAT WE HAVE TO UPDATE LOCATION AA OURSELVES WITH THE
 /"ISZ" INSTRUCTION. BY USING AUTO-INCREMENT REGISTERS
 /WE CAN GET AROUND THIS. *** SEE SOLUTION SUM2.

.TITLE SOLUTION USES INDIRECT ADDRESSING

00100

.ABS
.LOC 100

00100 777773

START

LAW =5
DAC COUNT#

/SET UP A COUNTER OF -5

00101 040135

00102 220125

NEXT

LAC* AA

/LOC AA CONTAINS THE ADD.
/OF A. LAC* AA GETS THE
/CONTENTS OF A IN AC.

00103 360126

TAD* BB

/ETC. FOR BB

00104 060127

DAC* CC

/ETC. FOR CC

00105 440125

ISZ AA

/INCREMENT THE INDIRECT

00106 440126

ISZ BB

/ADDRESSES OF AA, BB

00107 440127

ISZ CC

/AND CC.

00110 440135

ISZ COUNT

/SEE IF ALL DONE.

00111 600102

JMP NEXT

/NOT DONE. GET NEXT DATA.

00112 740040

HLT

/ALL DONE SO HALT.

00113 000001

A

1/ 2/ 3/ 4/ 5

00114 000002

00115 000003

00116 000004

00117 000005

00120 000002

B

2/ 3/ 4/ 5/ 6

00121 000003

00122 000004

00123 000005

00124 000006

00125 000113

AA

A

00126 000120

BB

B

00127 000130

CC

C

00130

C

.BLOCK 5

/RESERVE SPACE FOR TABLE C.

000100

.END START

SIZE=00136

NO ERROR LINES

.TITLE COMMENTARY ON SUM2

```

*****
*
* POINTS TO BE NOTED
*
*****

```

```

/*1. NOTE THAT IN THIS PROGRAM, DATA APPEARS BEFORE THE
/INSTRUCTIONS IN THE PROGRAM. THE IMPORTANT THING IS TO
/SEPARATE THE DATA FROM THE INSTRUCTIONS, BUT IT DOESN'T
/MATTER WHICH APPEARS FIRST.
/

```

```

/*2. NOTE THAT THE FIRST EXECUTABLE LOCATION IN A PRO-
/GRAM IS NOT ALWAYS IN THE FIRST LOCATION USED BY THE
/PROGRAM. THAT IS, THERE CAN BE A DIFFERENCE BETWEEN THE
/PROGRAM START ADDRESS AND THE PROGRAM LOAD ADDRESS.
/IN THIS CASE, THE PROGRAM START ADDRESS IS 117, WHEREAS
/THE PROGRAM LOAD ADDRESS IS 100.
/

```

```

/*3. THIS PROGRAM MAKES USE OF AUTO-INCREMENT ADDRESSING
/(WITH LOCATIONS 10, 11, AND 12). TO ACCESS SEQUENTIAL
/LOCATIONS IN TABLES A, B AND C RESPECTIVELY (LAC*10,
/TAD* 11, DAC*1 12).
/

```

```

THIS MEANS THAT WE WANT TO LOAD LOCATION 10 WITH A
/VALUE 1 LESS THAN THE START ADDRESS OF TABLE A,
/LOCATION 11 WITH A VALUE 1 LESS THAN THE START ADDRESS
/OF TABLE B AND LOCATION 12 WITH A VALUE 1 LESS THAN THE
/START ADDRESS OF TABLE C. NOTE THAT THE "LAC (A=1"
/AND "DAC 10" INSTRUCTIONS ARE USED (LOCATIONS 121 AND
/122) TO DO THIS.
/

```

```

THE USE OF "(A=1" REQUESTS THE ASSEMBLER TO SET UP
/A LOCATION CONTAINING THE VALUE "A=1". THE ASSEMBLER
/CAN EVALUATE EXPRESSIONS AND DOES SO MOVING FROM LEFT
/TO RIGHT...A=1 = 100-1 = 77. LOCATION 136 IS SET UP TO
/CONTAIN THE 77. THE INSTRUCTION "LAC (A=1" IS ASSEMBLED
/AS 200136 (LOAD THE ACCUMULATOR WITH THE CONTENT OF
/LOCATION 136, I.E., THE 77).
/

```

```

SIMILARLY "LAC (B=1" LOADS THE ACCUMULATOR WITH
/"104" AND "LAC (C=1" LOADS THE ACCUMULATOR WITH "111".
/

```

```

/*4. TO STORE THE "77" IN ABSOLUTE LOCATION 10 A
/"DAC 10" INSTRUCTION IS USED. WOULD WE BE LOADING THE
/"77" INTO ABSOLUTE LOCATION 10 IF THE ".LOC" STATEMENT
/READ ".LOC 20000" RATHER THAN ".LOC 100"?
/

```

.TITLE SOLUTION USES AUTO-INCREMENT ADDRESSING

00100

.ABS
.LOC 100

00100 000001
00101 000002
00102 000003
00103 000004
00104 000005
00105 000002
00106 000004
00107 000001
00110 000005
00111 000003
00112

A

11 21 31 41 5

B

21 41 11 51 3

C

.BLOCK 5

00117 777773
00120 040135

START

LAW -5
DAC COUNT#

/PUT -5 INTO COUNT

00121 200136
00122 040010

LAC (A-1
DAC 10

/THIS METHOD MAKES USE OF
/AUTO-INCREMENT REGISTERS

00123 200137
00124 040011

LAC (B-1
DAC 11

/10, 11 AND 12. SINCE
/AUTO-INCREMENT IS A PRE-

00125 200140
00126 040012

LAC (C-1
DAC 12

/INCREMENT, EACH OF THE
/REGISTERS MUST BE LOADED WITH

00127 220010
00130 360011
00131 060012
00132 440135

NXT

LAC* 10
TAD* 11
DAC* 12
ISZ COUNT

/ONE LESS THAN THE STARTING ADD

00133 600127
00134 740040

JMP NXT
HLT

/AN AUTO-INCREMENT REGISTER WIL
/WORK IN ANY PAGE OR BANK, THE
/INSTRUCTION TO LOAD LOCS 10, 11
/AND 12 SHOULD BE OF THE FORM

LAC (A-1
DAC* (10

.END START

00136 000077 *L
00137 000104 *L
00140 000111 *L

SIZE=00141

NO ERROR LINES

gilder andra
PAGES or Banks than 0

.TITLE COMMENTARY ON SUM3

```
*****  
*  
* POINTS TO BE NOTED *  
*  
*****
```

/*1. NOTE THE USE OF ".ABSP". THIS PROGRAM MAKES USE OF
/INDEXED ADDRESSING AND, THEREFORE, MUST BE IN PAGE MODE

/*2. NOTE THE USE OF THE "DBA" INSTRUCTION. ALTHOUGH
/A "DBA" INSTRUCTION IS CONTAINED IN THE PAGE MODE
/VERSION OF THE ABSOLUTE LOADER (WHICH IS OUTPUT ON THE
/PAPER TAPE IN FRONT OF YOUR ASSEMBLED PROGRAM), IT IS
/STILL A GOOD IDEA TO INCLUDE A "DBA" IN YOUR PROGRAM
/IN CASE THE PROGRAM IS RESTARTED FROM THE CONSOLE
/WITHOUT RELOADING IT. DON'T RELY ON THE CONSOLE
/BANK-PAGE MODE SWITCH.

/*3. NOTE THE USE OF THE ".DEC" AND ".OCT" ASSEMBLER
/DIRECTIVES. COMPARE THE VALUES GENERATED FOR THE
/NUMBERS STATED WHILE UNDER DECIMAL RADIX (BASE 10).
/WITH THOSE GENERATED FOR THE NUMBERS STATED WHILE
/UNDER OCTAL RADIX (BASE 8).

/*4. NOTE THE ".END BEGIN=1" STATEMENT. A ".END"
/STATEMENT IS USED TO SPECIFY THE PHYSICAL END OF
/A PROGRAM. A "START ADDRESS" OR "TRANSFER ADDRESS"
/(THE LOCATION AT WHICH WE WANT THE LOADER TO START THE
/PROGRAM) MAY ALSO BE SPECIFIED IN THE ".END" STATEMENT.
/THIS IS DONE BY FOLLOWING THE ".END" WITH A SPACE OR
/TAB AND THEN GIVING AN EXPRESSION WHICH MAY BE EVALUA-
/TED BY THE ASSEMBLER. IN THE PREVIOUS EXAMPLES,
/WHEN A TRANSFER ADDRESS WAS GIVEN IT WAS A VERY
/SIMPLE EXPRESSION--THE LABEL USED ON THE FIRST
/EXECUTABLE LOCATION. BUT THE ASSEMBLER CAN HANDLE
/LONGER EXPRESSIONS. IN THIS CASE, THE SYMBOL BEGIN
/THE VALUE 10001, THEREFORE, BEGIN=1=10000,
/THUS, IN THIS CASE, ".END BEGIN=1" HAS THE SAME
/EFFECT AS ".END 10000".

/*5. NOTE THE ERROR DIAGNOSTIC GIVEN ON THE STATEMENT
/TO BE ASSEMBLED INTO LOCATION 10025. THE "N" INDICATES
/AN ERROR IN NUMBER USAGE. WHAT IS THE ERROR?
/HOW WAS IT HANDLED BY THE ASSEMBLER (COMPARE LOCATIONS
/10020 AND 10025)?

PAGE	2	SUM3	SRC	ROUTINE USES INDEXED ADDRESSING (8 LIMIT REGISTER)
				.TITLE ROUTINE USES INDEXED ADDRESSING (8 LIMIT REGISTER)
10000				.ABSP .LOC 10000
10000	707762			DBA /ENTER PAGE MODE
10001	200033	BEGIN		LAC (5
10002	722000			PAL /PUT 5 IN LIMIT REG.
10003	735000			CLX /XR=0
10004	210014	AA		LAC A,X /ADDRESS OF A + C(XR).
10005	350021			TAD B,X /ETC,
10006	050026			DAC C,X /ETC.
10007	725001			AXS 1 /INCREMENT AND TEST XR
10010	600004			JMP AA /XR & LR
10011	740040			HLT /XR=LR. WE'RE DONE!
10012	000000			CAL /CHECK RESULTS VIA CONSOLE.
10013	000015			IS /RETURN TO MONITOR
				/BY HITTING CONSOLE CONTINUE.
10014	000002	A		.DEC
10015	000005			2/ 5/ 12/ 27/ 58
10016	000014			
10017	000033			
10020	000072			
10021	000002	B		.OCT
10022	000005			2/ 5/ 12/ 27/ 58
10023	000012			
10024	000027			
10025	000072			
10026		C		.BLOCK 5
10033	010000			.END BEGIN=1
	000005	*L		1 ERROR LINES
		SIZE=10034		

)

)

)

)

)

)

)

)

/PROGRAM TO ADD TWO NUMBERS, USING TAD, AND TEST
/THE RESULT FOR ARITHMETIC OVER FLOW.

```

START   CLA!CLL           /CLEAR THE LINK AND THE AC.
        TAD A             /GET THE FIRST # IN THE AC.
        AND MASK         /GET RID OF ALL BITS EXCEPT THE SIGN.
        TAD B             /ADD THE 2ND NO. TO BIT 0 OF A.
        SZL              /SKIP IF ZERO LINK
        JMP NEGNEG       /NON ZERO LINK INDICATES BOTH A AND B NEG.
                               /BECAUSE ONES IN BOTH SIGN POSITIONS IS
                               /THE ONLY WAY THE LINK CAN BE SET SINCE ONLY
                               /THE SIGN BIT OF THE 2ND # IS ADDED
        RAL              /GET RESULT SIGN INTO LINK.
        SZL              /SKIP IF ZERO LINK
        JMP POSNEG       /LINK=1, MUST BE +-.
        JMP POSPOS       /LINK=0, MUST BE ++.
POSNEG  LAC A             /SINCE ONE VALUE IS + ADD
        TAD B             /THE OTHER IS -. CORRECT ADDITION IS ASSURED.
        DAC SUM          /NO TEST IS NECESSARY, SO STORE THE SUM
        HLT              /AND HALT.
POSPOS  LAC A             /IF THE SUM OF TWO POSITIVE NUMBERS GIVES
        TAD B             /A NEGATIVE RESULT, ARITHMETIC OVERFLOW HAS
        SPA              /OCCURED, TEST THE AC FOR POSITIVE VALUE.
        JMP POSERR       /IF NEGATIVE, GO TO ERROR ROUTINE.
        DAC SUM          /POSITIVE RESULT SO SUM OK, STORE RESULT
        HLT              /AND HALT.
NEGNEG  LAC A             /IF THE SUM OF TWO NEGATIVE NUMBERS GIVES
        TAD BQ           /A POSITIVE RESULT, ARITHMETIC OVERFLOW HAS
        SMA              /OCCURED, TEST THE AC FOR NEGATIVE VALUE.
        JMP NEGERR       /IF POSITIVE, GO TO ERROR ROUTINE.
        DAC SUM          /NEGATIVE RESULT SO SUM OK, STORE RESULT
        HLT              /AND HALT.
MASK    400000

```

/ANOTHER, SHORTER SOLUTION FOLLOWS:

```

        LAC A             /IF THE OR'ED SUM OF THE SIGN
        XOR B             /BITS IS A 1, THE SIGNS WERE
        SMA              /DIFFERENT AND THE SUM MUST BE CORRECT
        JMP LIKE         /SIGNS WERE THE SAME
        LAC A
        TAD B
        DAC SUM
        HLT
LIKE    LAC A             /PERFORM THE ADDITION, AND THEN
        TAD B             /CHECK THE SUM FOR A CORRECT SIGN
        DAC SUM          /
        AND (400000)     /GET SIGN OF RESULT. IF OR'ED SUM
        XOR B             /GIVES A SIGN OF 0, THE SIGN OF THE
                               /RESULT IS THE SAME AS THE SIGN OF B.
        SPA              /SINCE A & B WERE SAME SIGN, RESULT IS OK
        JMP ERROR        /SIGN CHANGED, OVERFLOW OCCURED
        HLT

```

/A THIRD SOLUTION MAKING USE OF ADD IS MUCH SIMPLER.

```

        CLL              /CLEAR LINK
        LAC A             /ADD THE TWO NUMBERS
        ADD B             /IF EITHER NO. IS NEG., IT MUST BE 1'S COMP.
        SZL              /IF LINK IS SET, OVERFLOW
        JMP ERROR        /OCCURED, GO TO ERROR
        DAC SUM          /ADDITION OK IF LINK=0
        HLT

```


)

)

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)

ADDRESSING

References: Volume 1 Processor Handbook 4-3
System Reference Manual 8-1

Memory Reference Instructions specify locations to be operated on by the Central Processing Unit. The CPU computes the actual (effective) address of the location referred to by combining bits from the instruction itself and also from the PC at the time the instruction is being executed. Various addressing modes require further computations with pointer words and the index register.

In the following illustrations:

PC: refers to the contents of the Program Counter at the time the MRI is being executed.

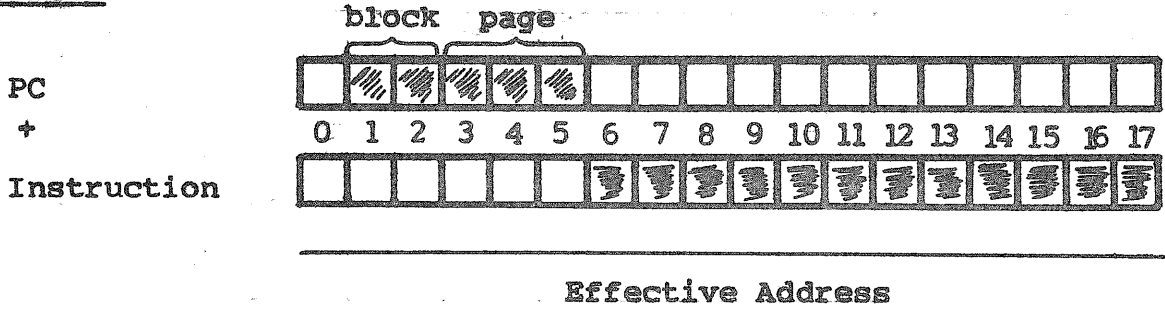
Instruction: refers to the contents of the location being executed.

XR: refers to the contents of the Index Register.

Pointer Word: refers to the contents of the location designated as a pointer word in an indirect reference.

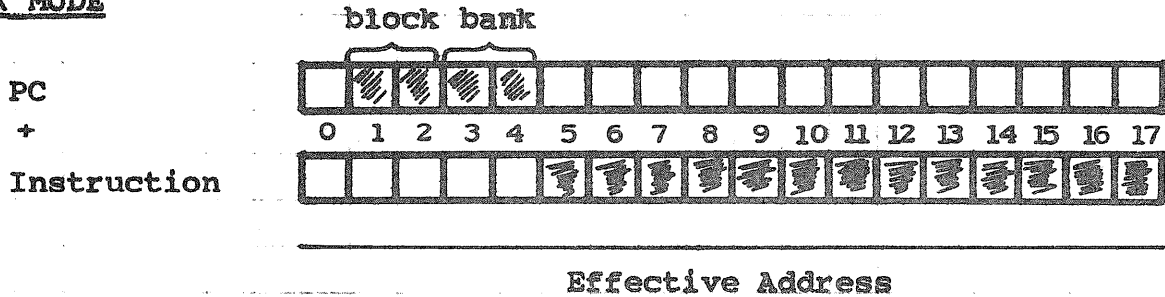
DIRECT ADDRESSING

PAGE MODE



NOTE that the effective address is in the same page as the instruction.

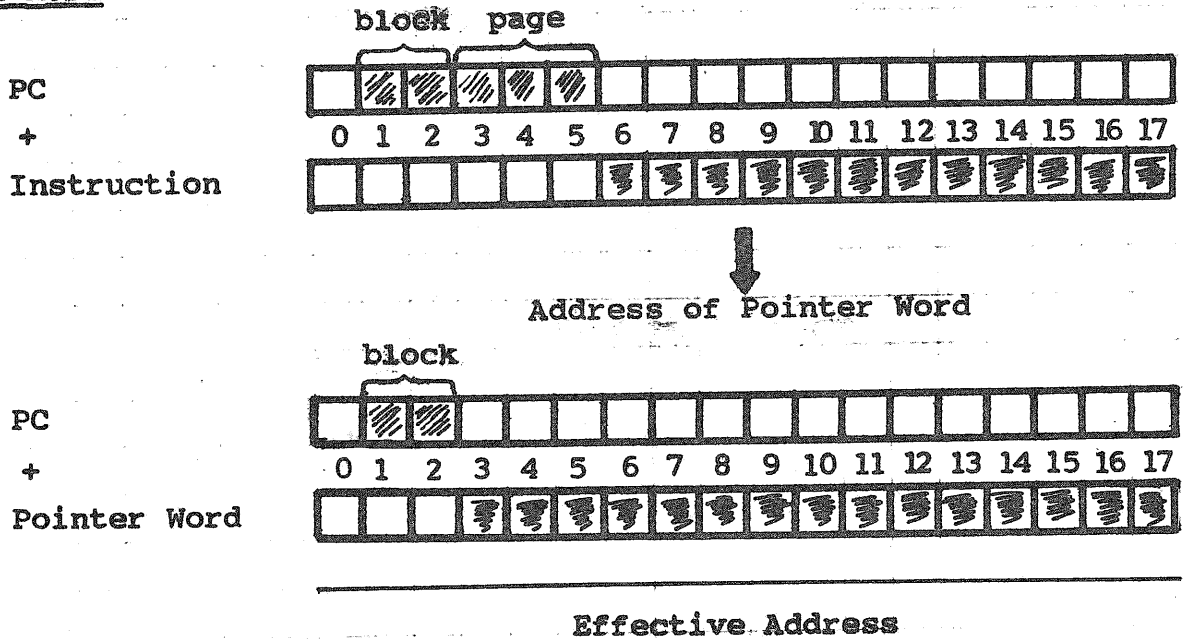
BANK MODE



NOTE that the effective address is in the same bank as the instruction.

INDIRECT ADDRESSING

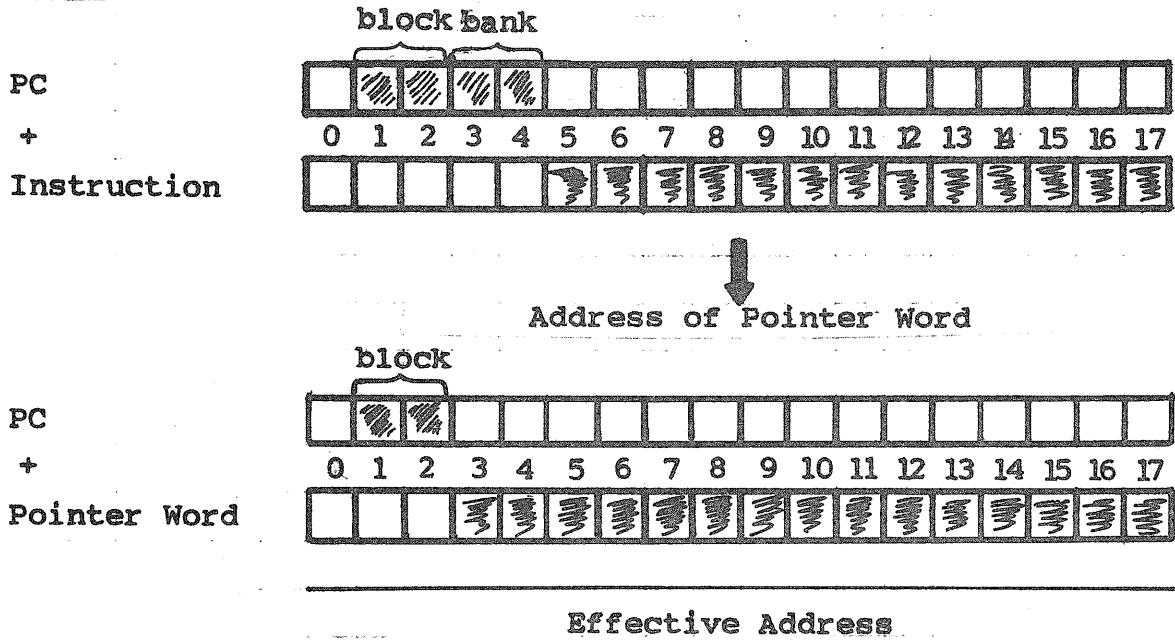
PAGE MODE



NOTE that the pointer word is in the same page as the instruction.
 NOTE that the effective address is in the same block (32K) as the instruction.

INDIRECT ADDRESSING

BANK MODE

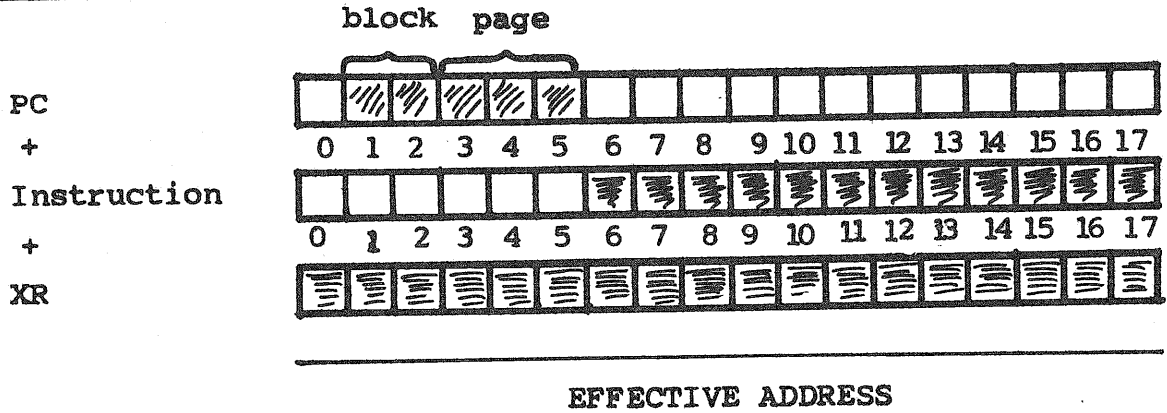


NOTE that the pointer word is in the same bank as the instruction.

NOTE that the effective address is in the same block (32K) as the instruction.

INDEXED ADDRESSING

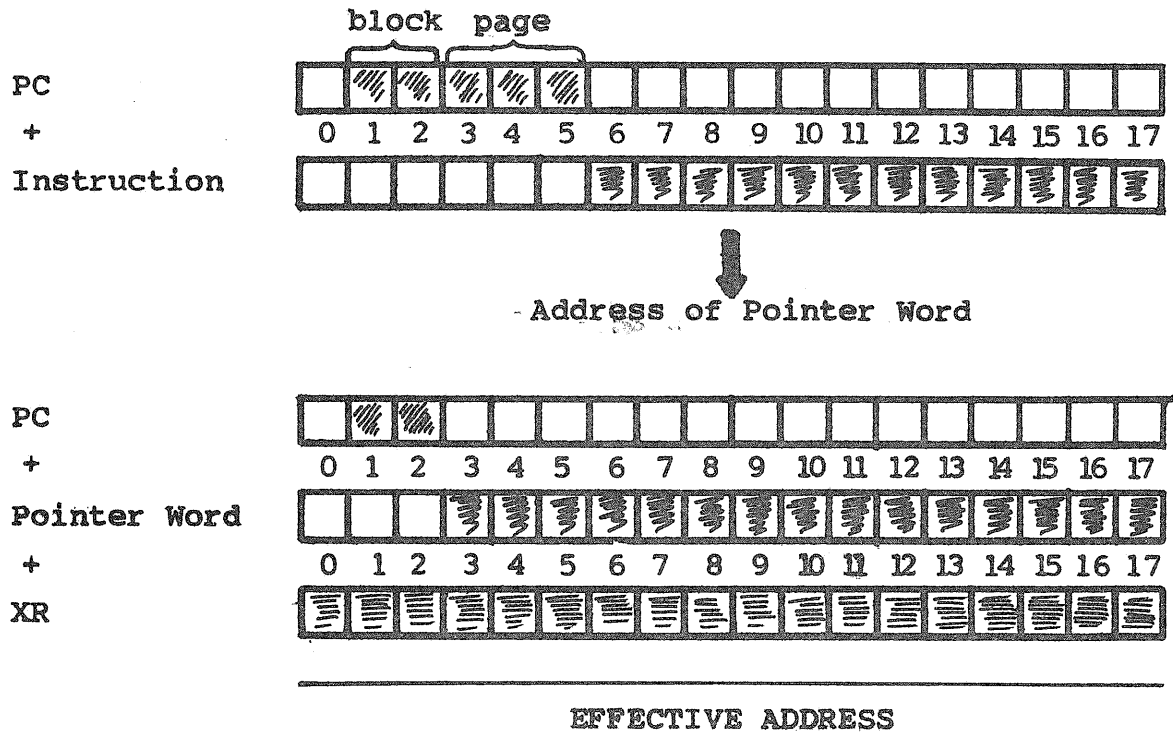
PAGE MODE ONLY



NOTE that the effective address may be anywhere in 128K.
 Care must be taken not to address non-existent memory
 (this includes negative addresses).

INDIRECT INDEXED ADDRESSING

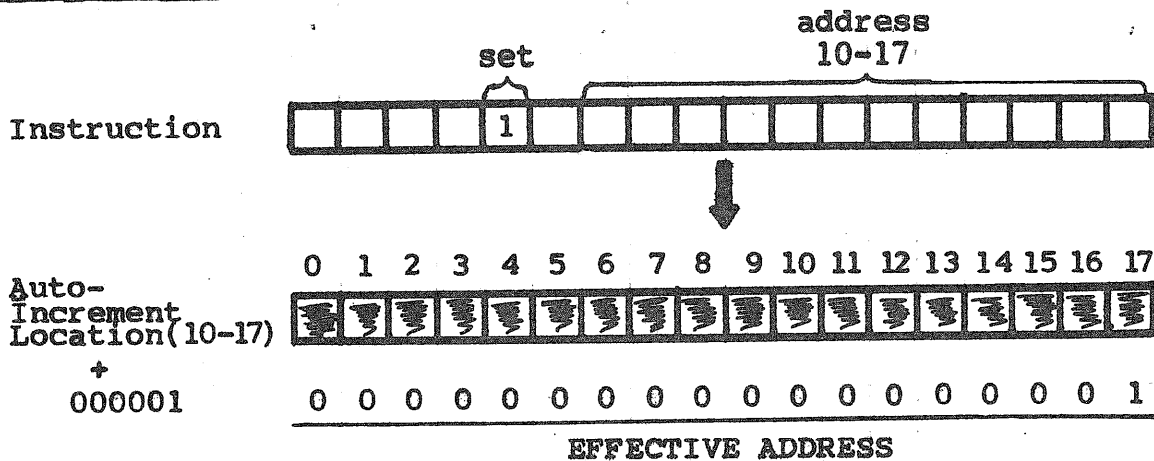
PAGE MODE ONLY



NOTE that the effective address may be anywhere in 128K.
 Care must be taken not to address non-existent memory
 (this includes negative addresses).

AUTOINCREMENT ADDRESSING

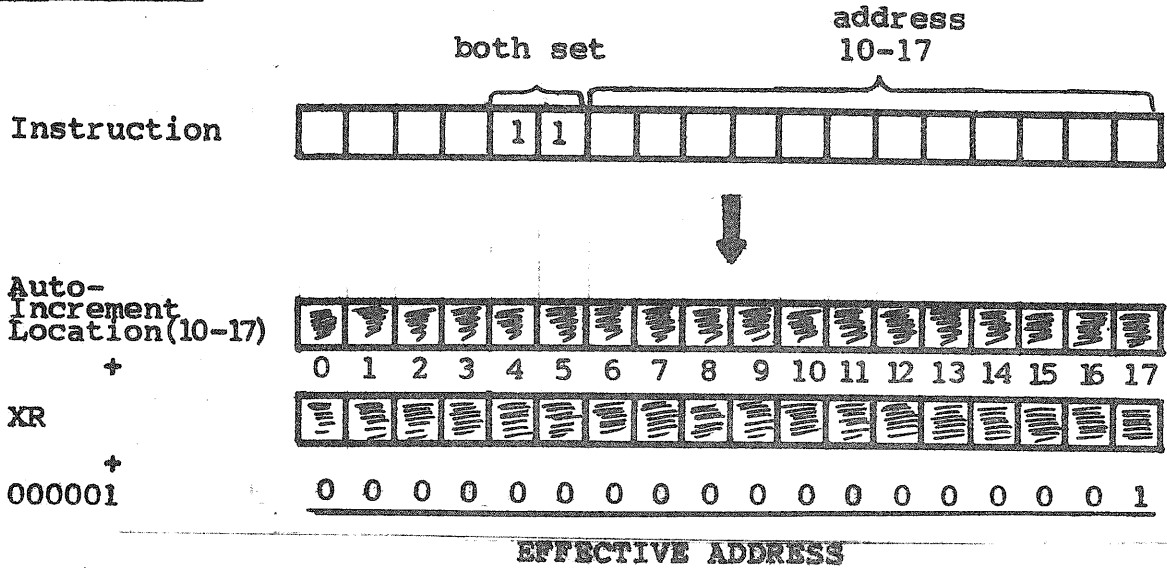
PAGE MODE and BANK MODE



NOTE that the effective address may be anywhere in 128K.
 Care must be taken not to address non-existent memory
 (this includes negative addresses).

INDEXED AUTOINCREMENT ADDRESSING

PAGE MODE ONLY



NOTE that the effective address may be anywhere in 128K.
 Care must be taken not to address non-existent memory
 (this includes negative addresses).

SUBROUTINES

What: A section of code, usually performing one task, that may be called from various points of the main program.

How: The JMS instruction is used to enter subroutines. Recall that upon a JMS

- a) The updated PC is stored at the address specified in the JMS.
- b) The (specified address) +1 is now placed in the PC so that execution is picked up at the second location of the subroutine.

<u>ex.</u>	100	Before	JMS	100	After	JMS
PC →	101	JMS	MOVE	101	JMS	MOVE
		⋮		⋮	⋮	
		200 MOVE	∅	200 MOVE	101	
				PC → 201		

c) Return to the calling program is effected via a "JMP*" on the first location of the subroutine.

```

MOVE    ∅
        ⋮
        JMP* MOVE
    
```

Passing Arguments - very often a subroutine has to process data contained in the calling routine. In order to do this the calling routine must pass this data or the address(es) of the data to the subroutine. This is called "passing arguments."

There are a number of ways this can be done. Among them are:

- | | | |
|------------------------|---|------------------------------------|
| 1) AC | } | passing arguments in CPU registers |
| 2) MQ, XR, LR | | |
| 3) trailing arguments: | | |
| | | JMS SUB |
| | | Arg 1 |
| | | Arg 2 |
| | | Arg 3 |

4) Setting up a list of arguments and passing the address via (1), (2), or (3).

LAC	(LIST ADDRESS	or	JMS	SUB	etc.
JMC	SUB		LIST	ADDRESS	
	⋮				
	⋮				
	⋮				

LIST ADDRESS	ARG1
	ARG2
	⋮
	ARGN

PC15 High-Speed Paper-Tape Reader Punch

1.1 INTRODUCTION

The PC15 High-Speed Paper Tape Reader/Punch is used to input perforated paper-tape programs into core memory, or to punch core memory programs or data on paper tape. Information is punched on 8-channel fanfolded paper tape in the form of 6- or 8-bit characters at a maximum rate of 50 characters/second. Information is read at a maximum rate of 300 characters/second. The PC15 consists of a PC05 Paper Tape Reader/Punch with interface and control logic for using the reader/punch with a PDP-15.

1.2 PAPER-TAPE READER

1.2.1 Characteristics and Capabilities

Data can be read from tape and transferred to the PDP-15, using the computer hardware readin logic or using program-controlled transfers. For hardware readin operation, the hardware readin logic supplies inputs for selecting the operating mode, starting tape motion, and implementing transfers. For program-controlled transfers, the computer issues input/output transfer (IOT) instructions that select the operating mode, advance the tape, and implement the transfer. To maintain a maximum rate of 300 characters/second, a new select IOT must be issued within 1.67 ms of the last reader flag. If not, the reader operates start-stop and reads characters at a 25 character/second rate. The requirements for maximum character rate are described in detail in Programming Considerations, Paragraph 1.4.1.

The reader interfaces with the automatic priority interrupt (API) facility, the program interrupt facility, and the input/output skip chain. For API operation, the reader is assigned API level 2; a unique entry address of 50_g is assigned to its service routine.

The reader contains a no-tape sensor and flag (character ready for transfer) circuits. If a no-tape condition is detected, the reader flag is set, and a program interrupt is initiated whenever a reader select IOT is given. The states of the reader flag, the reader API 2 level, PI request and skip request

devices are displayed on an indicator panel at the top of Cabinet H963E (Bay 1R). In addition, this panel displays the reader buffer contents and the I/O address (API unique entry address). These items and the reader controls are described in Controls and Indicators, Paragraph 1.2.3.

Reader mechanical facilities include a right-hand bin for supply for tape being read, a left-hand bin for receiving the tape, and a feed-through mechanism to control passage of the tape into the receiving bin. A snap-action retainer on the feed-through mechanism facilitates simple loading of the tape.

1.2.2 Operating Modes

The PC15 reader operates in either an alphanumeric or binary mode. For program-controlled transfers, the operating mode is selected by IOT instructions. For hardware readin operation, control logic in the reader automatically selects the binary mode.

When alphanumeric mode is selected, one 8-bit character (in ASCII code) is read and transferred to the PDP-15 accumulator. In the binary mode, the reader reads three 6-bit characters (three frames with channels 7 and 8 ignored) from tape and assembles them into an 18-bit word for transfer to the accumulator.

1.2.3 Controls and Indicators

Two front panel controls are provided for the PC15 Paper-Tape Reader: ON LINE/OFF LINE and FEED. The ON LINE position places the reader under computer control. The OFF LINE position, which is used for loading paper tape, raises an out-of-tape flag and places the reader under local control. The indicators associated with reader operation are located on an indicator panel at the top of cabinet H963E (Bay 1R). Table 1-1 lists the indicators and their functions.

Table 1-1
Indicators Associated with Paper-Tape Reader

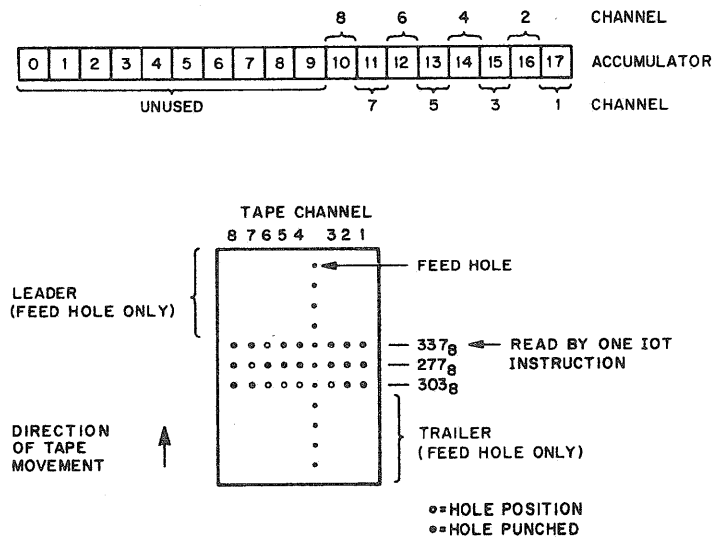
Indicator	Function
READER BUFFER 00-17 API 2 RDR	Indicates the contents of the paper-tape reader buffer. Denotes API level 2 is active as the result of a reader interrupt.
I/O ADDRESS	Indicates the unique trap address associated with I/O devices; address 50g for paper-tape reader.
RDR FLG	Denotes information has been read from tape and is available for transfer from reader buffer.

Table 1-1 (Cont)
Indicators Associated with Paper-Tape Reader

Indicator	Function
PI RQ	Denotes one of the I/O devices (including paper-tape reader) handled by the BA15 Peripheral Expander has generated an interrupt request.
SKIP RQ	Denotes one of the I/O devices (including paper-tape reader) handled by BA15 has responded to a skip IOT instruction.

1.2.4 Tape Formats

The format of the perforated paper tapes for the alphanumeric (ASCII usage) mode is shown in Figure 1-1. In addition, tape channels are related to the PDP-15 accumulator stages. The leader and trailer portions of the tape are used to introduce or conclude a paper-tape program. Only the feed hole is punched for the leader/trailer portions. Note that each character is read by one IOT instruction.



15-0232

Figure 1-1 Tape Format and Accumulator Bits (Alphanumeric Mode)

The paper-tape format for binary mode using hardware readin (HRI) is shown in Figure 1-2 as well as the relationship of accumulator stages for the 18-bit word. Note that only the feed hole is perforated for the leader/trailer portion and that channel 8 is always punched in the program portion of the tape. Any character without hole 8 punched will be ignored. Channel 7 punched in the last character indicates the last 18-bit instruction is to be executed by the computer. This instruction can halt machine operation or can transfer machine control to another part of the program. When using this format, channel 7 must be punched using the alphanumeric mode.

1.2.5 Instructions

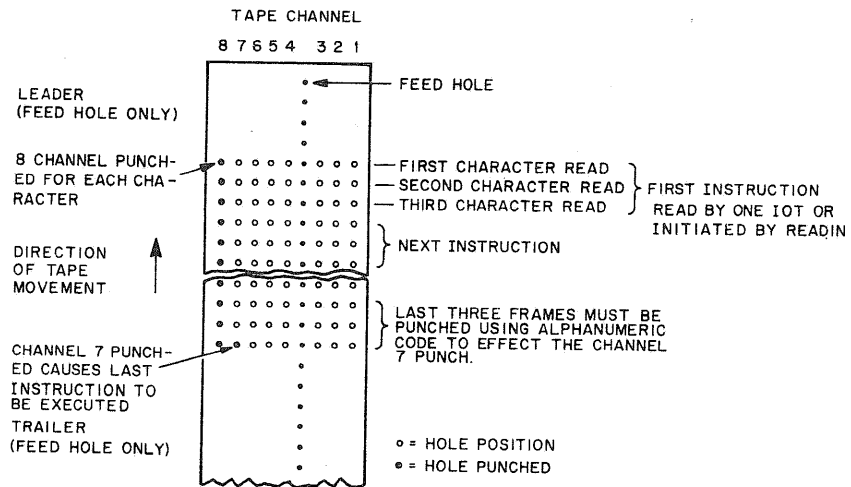
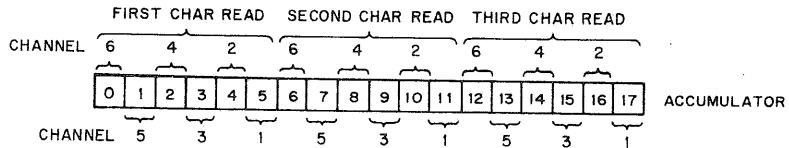
The PDP-15 IOT instructions used for program-controlled loading of paper-tape data are listed below. Refer to Volume 1 of this handbook for IOT instruction format.

<u>Mnemonic</u>	<u>Octal Code</u>	<u>Operation Performed</u>
RSF	700101	Skip next instruction if reader flag is a 1.
RCF	700102	Clear reader flag. Read reader buffer, inclusively OR contents of reader buffer with AC, and deposit result in AC.
RRB	700112	Read reader buffer and clear reader flag. Clear AC and transfer contents of reader buffer to AC.
RSA	700104	Select alphanumeric mode and place one 8-bit character in reader buffer. Clear flag before character is read from tape. Set reader flag to 1 when transfer to reader buffer is complete.
RSB	700144	Select binary modes. Assemble three 6-bit characters in reader buffer. Clear reader flag during assembly and set flag when assembly is complete.

The paper-tape reader responds to an input/output read status (IORS) instruction by supplying the status of its device flags and no-tape flags to the accumulator. The reader device flag (reader interrupt) interfaces with bit 01 of the accumulator. The reader no-tape flag interfaces with bits 08 of the accumulator.

1.2.6 Functional Description

The PC15 reader consists of an electromechanical tape feed system, a light source and photo cells for sensing tape perforations, a buffer register for storing and assembling data, and control logic for computer interface, tape advance, and transfer operations. These circuits can be used with the PDP-15 hardware readin logic, or can be used for program-controlled transfers, as described in the following paragraphs.



15-0233

Figure 1-2 HRI Tape Format and Accumulator Bits (Binary Mode)

1.2.6.1 Hardware Readin Operation - The PC15 reader can be used with PDP-15 hardware readin logic to load programs from paper tape at a rate of 300 characters/second. For this operation, the desired tape is installed in the high-speed reader, and the program loading address is selected, using the console ADDRESS switches. The console RESET key is then pressed to initialize the computer and paper-tape reader. A readin operation is started by pressing the READIN key on the console.

With this key action, a Readin (RI) condition is stored in the reader, and the binary mode is selected. The reader then advances the tape, reads three characters from tape, assembles them into an 18-bit word in the reader buffer, and signals the hardware readin logic with a program interrupt. The hardware readin logic, in turn, transfers the 18-bit word to the accumulator under I/O processor and computer timing. The word is subsequently loaded into core memory by forcing a DAC instruction. The first 18-bit word is stored at the address specified by the console ADDRESS switches. Subsequent 18-bit words are stored in sequential memory locations.

The readin operation continues until a perforated hole 7 is detected. This condition is inserted in the last character of the last 18-bit instruction. When this condition is detected, the reader supplies the hardware readin logic with a skip request. As a result, the hardware readin logic causes the last instruction to be loaded into the Memory Input register for execution. This instruction can halt machine operation (HALT) or can transfer program control to another part of the program (JMP). When using the readin feature with the MP15 Memory Parity option, the last instruction on the paper tape (which will be executed by the processor) will not be written into the next sequential memory location. That location, however, will be loaded with data that may contain wrong parity. Therefore, that location should be re-stored by the program before an attempt is made to read from it. Otherwise, a parity error will occur.

1.2.6.2 Program-Controlled Operation - The PC15 reader operates in the binary or alphanumeric mode depending on the select IOT instructions issued by the computer. On decoding a reader select alphanumeric (RSA) mode IOT (700104g), the reader advances the tape one character, loads this character into the reader buffer, and sets the reader device flag. The reader then signals the computer that data are available by providing a reader interrupt to the API or PI, or by responding to an RSF IOT instruction. If the API facility is being used, program control is transferred to the reader service routine where the computer services the request, and an RCF (700102g) or RRB (700112g) instruction is issued. If the API facility is not being used, the computer issues an RSF instruction, and the reader returns a skip request whenever its flag is set. The skip request causes the next instruction (normally a JMP .-1 in wait loops) to be skipped so that the character can be transferred to the accumulator by issuing an RCF or RRB instruction. The RCF or RRB instruction transfers the reader buffer character to the I/O bus and loads it into the least significant bits (10 through 17 for 8-bit alphanumeric character) of the accumulator. The character is subsequently stored in a core memory location designated by the program. The read reader buffer (RRB) instruction also clears the reader flag for the next read operation.

For binary mode operation, the computer issues a reader select binary (RSB) mode instruction (octal 700144). On decoding this instruction, the reader clears its device flag, advances the tape three

characters, reads these characters from tape, and assembles them into an 18-bit word in the reader buffer. The reader also counts the number of characters with hole 8 punched read from tape and, when a count of three is reached, generates an interrupt request. The control functions for transfer of the 18-bit word to the accumulator is the same as that described for the alphanumeric mode.

1.3 PAPER-TAPE PUNCH

1.3.1 Characteristics and Capabilities

The PC15 paper-tape punch consists of a tape feed system, a mechanical punch assembly, a buffer register, and control logic for mode selection and activation of the tape feed and punch mechanism. Tape advance, mode selection, and transfer of information to the punch are controlled by IOT instructions. Tape is perforated at a rate of 50 characters/second. When the punch is selected by an IOT instruction, data from the PDP-15 accumulator (AC10-AC17) are transferred to the punch buffer. Then, without further inputs, a character is perforated on tape.

The punch contains a device flag that denotes punch status for transfers. This device flag interfaces with the PI facility and I/O skip chain. The status of the punch flag is displayed on an indicator panel at the top of Cabinet H963E (Bay 1R). An out-of-tape switch is located on the punch mechanism. This switch initiates action that stops punch operations when approximately one inch of unpunched tape remains.

Power for the punch operation is available whenever the PDP-15 power is on. The punch runs when selected by an IOT instruction or when the FEED switch is pressed.

Punch mechanical features include a magazine for unpunched tape and a container for tape chad. Both are accessible when the reader-punch drawer is extended from the cabinet.

1.3.2 Operating Modes

The PC15 Punch operates in the alphanumeric or binary mode as designated by IOT select instructions. One of these instructions is required for each character punched for mode change. In the alphanumeric mode, an 8-bit character (in ASCII or modified ASCII code) is punched for each accumulator transfer to the punch. For the binary mode, one 6-bit data character is perforated for each accumulator transfer. Hole 8 is always punched, and hole 7 is never punched. Three of these characters, however, form one computer word for readin operations.

1.3.3 Controls and Indicators

The PC15 Punch has a front panel FEED control. This control is used to advance the tape from the punch as required for leader or trailer. The punch also has one indicator (PUN FLG) directly associated with its operation. This indicator, located on an indicator panel at the top of Cabinet H963E (Bay 1R), indicates the status of the device flag and, shows that the punch is available for a punch operation when lit. The punch also shares the PI RQ and SKIP RQ indicators on this panel with other I/O devices.

1.3.4 Tape Formats

Tape formats are shown in Figures 1-1 and 1-2.

1.3.5 Instructions

The PDP-15 IOT instructions used for punching of paper tape under program control are listed below. Refer to Volume 1 of this handbook for IOT instruction format.

<u>Mnemonic</u>	<u>Octal Code</u>	<u>Operation Performed</u>
PSF	700201	Skip next instruction if punch flag is a 1.
PCF	700202	Clear punch flag and punch buffer.
PSA	700204	Select alphanumeric mode and punch one character. Set punch flag when punch is complete.
PSB	700244	Select binary mode and punch one 6-bit character. Set punch flag when punch is complete.

The punch responds to the IORS instruction (Volume 1, Paragraph 3.7.1) by supplying the status of its device flag and no-tape flag to the accumulator. The device flag interfaces with bit 02 of the accumulator, and the no-tape flag interfaces with bit 09.

1.3.6 Functional Description

The PC15 Punch operates in the alphanumeric or binary mode, depending on whether a PSA or PSB instruction is issued. When one of these instructions is decoded, information is loaded into the punch buffer from bits 10 through 17 of the accumulator and is punched onto tape. During the interval the punch operation is in progress, the punch flag is cleared to indicate the punch is busy. When the punch operation is complete, the punch flag is set to 1 to indicate it can accept another input character.

The operating sequence for punch operations normally begins with a PSF instruction to test the device flag. If the device flag is 1, a skip request is returned to the computer, and the computer issues a PCF instruction. This instruction clears the device flag and the punch buffer. The computer then issues a PSA or PSB instruction. On decoding a PSA instruction, the reader loads the accumulator input into its buffer, advances the tape, and punches one character. For the alphanumeric mode channel 8 is punched as a function of bit AC10. For the alphanumeric mode channel 7 is perforated as a function of bit AC11. After the character is punched, the reader sets its device flag, and the process is repeated. This operation, performed by the PCF and PSA instructions, can be combined by microprogramming the two instructions to form octal 700206.

The same principles are used for punching a binary character; however, a PSB instruction is used in place of the PSA instruction. On decoding a PSB, the punch perforates channel 8 and inhibits the punching of channel 7. The remaining six channels are punched as a function of AC12 through AC17, and represent one 6-bit character of a computer word.

1.4 PROGRAMMING CONSIDERATIONS

1.4.1 High-Speed Paper-Tape Reader

To use the reader at the transfer rate of 300 cps, a select IOT (RSA or RSB) must be issued within 1.67 ms after each flag. This action is required because a 40 ms reader stop delay is present. When this delay is activated, it overrides the select IOT input and subsequently stops the tape. Thus, if a new select IOT is not received within 1.67 ms of the setting of the flag, the reader operates start-stop and reads characters at 25 cps rate. No data is lost.

The RSA (octal 700104) and RCF (octal 700102) can be microprogrammed to form an octal 700106 instruction. This instruction reads the character, transfers the character to the accumulator, and advances the tape in one operation. An RSF (octal 700101) and RRB (700112) cannot be microprogrammed.

1.4.2 High-Speed Paper-Tape Punch

Channel 7 can be punched using only the alphanumeric mode. Therefore, when punching the last character of a tape for hardware readin operation, the last character must be punched in the alphanumeric mode.

The PCF instruction can be microprogrammed with a PSA or PSB instruction to form octal 700206 or 700246. This instruction clears the punch flag and buffer, selects the applicable mode, loads the

punch buffer, advances the tape, and perforates the character on tape. After completing the punching, the punch flag is set to denote the punch can accept another character. Microprogramming the PCF and PSF instructions is not allowed.

1.5 PROGRAMMING EXAMPLES

1.5.1 Paper-Tape Reader/Punch Handlers

All PDP-15 Systems are supplied with standard I/O device handler subroutines for the paper-tape reader/punch hardware. For PDP-15/10 Systems with 4K core, the COMPACT software includes paper-tape handler routines such as PTLIST and PTDUP. The Basic I/O Monitor, supplied with PDP-15/10E Systems with 8K core or greater, include standard I/O device handlers for the high-speed paper-tape reader and punch. These standard device handlers operate in systems with or without API and are upward compatible with all other monitors on the PDP-15/20 Software System. Complete instructions on use of standard paper-tape reader and punch handlers and their modification for special applications are provided in the PDP-15/10 Software System Manual, DEC-15-GRIA-D.

1.5.2 Paper-Tape Reader Programming Example

The following subroutine illustrates the use of programmed IOT instructions to read a group of binary words from paper tape. Twenty-five 18-bit words are read and stored in a table starting at ADDRESS.

NOTE

This example is for instructional purposes only and is not to be considered a complete, fully tested software system segment.

SUBRTE	0		
	LAW	-31	/25 DECIMAL WORDS
	DAC	WDCNT	
	LAC	(ADDRESS	
	PAX		
READLP	IORS		
	AND	(1000	/IS THE PAPER TAPE READER EMPTY?
	SZA		/YES IF NON-ZERO.
	JMP*	SUBRTE	/EXIT...ITS EMPTY
	RSB		/NO. START READING A WORD.
	RSF		
	JMP	.-1	/WAIT FOR IT.
	RRB		/GET IT FROM HARDWARE BUFFER.
	DAC	0,X	
	AXR	1	/POINT TO NEXT LOC AT ADDR.
	ISZ	WDCNT	/HAVE 25 WORDS BEEN READ?
	JMP	READLP	/NO...CONTINUE LOOPING.
	JMP*	SUBRTE	/YES. EXIT.

1.3.3 Paper-Tape Punch Programming Example

The following subroutines illustrate some paper-tape punch programming considerations. Their purpose is to unpack successive 6-bit ASCII characters from a table, convert them to 7-bit ASCII, and punch them on paper tape. The starting address of the table is placed in a location named ADDRESS. The number of words in the table is placed in WORDCNT. After these parameters have been deposited, the subroutines are entered by a JMS to PNCHOUT.

NOTE

This example is for instructional purposes only and is not to be considered a complete, fully tested software system segment.

PNCHOUT	0		
	LAC	WORDCNT	/THIS INITIALIZATION
	TCA		/ROUTINE STORES 2'S
	DAC	WORDCNT	/COMPLEMENT WORDCNT
	CLX		/AND CLEARS XR.
NXTWORD	LAW	-3	/SET UP A COUNTER FOR
	DAC	COUNT	/3 CHARACTERS.
	LAC	ADDRESS,X	/USE XR TO GET EACH WORD.
	RAL		/AC HOLDS 3 6-BIT ASCII
			/CHARS. ROTATE INTO LINK.
NXTCHAR	RTL		/ROTATE WORD 6 PLACES
	RTL		/THRU LINK. THE NEXT
	RTL		/6-BIT CHAR. IS IN AC12-17.
	DAC	SAVEAC	/SAVE REMAINING CHARS.
	AND	(77)	/THIS ROUTINE CONVERTS
	TAD	(40)	/THE 6-BIT ASCII IN AC12-17
	AND	(77)	/TO 7-BIT ASCII IN
	TAD	(40)	/AC11-17.
	JMS	PPCHAR	/READY TO PUNCH CHAR.
	LAC	SAVEAC	/RESTORE SHIFTED AC.
	ISZ	COUNT	/LAST CHARACTER?
	JMP	NXTCHAR	/NO. DO NEXT CHARACTER.
	AXR	1	/POINT TO NEXT WORD.
	ISZ	WORDCNT	/LAST WORD?
	JMP*	PPASCII	/NO. DO NEXT WORD.
	JMP*	PNCHOUT	/YES. RETURN TO PROGRAM.
PPCHAR	0		
	DAC	STORE	/SAVE CHAR. FOR 'NO TAPE'
			/TEST.
	IORS		/LOAD PUNCH STATUS INTO AC.
	AND	(400)	/TEST NO PUNCH TAPE BIT.
	SZA		/SKIP IF TAPE OK.
	JMP	EOT	/GO TO END OF TAPE RTE.

```

LAC   STORE           /LOAD AC WITH CHARACTERS
PSA
PSF
JMP   .-1             /SELECT ALPHA MODE & PUNCH
PCF
JMP*  PPCHAR         /WAIT FOR FLAG.
                          /RETURN TO SUBPROGRAM.

```

1.5.4 Programming With API or PI

The standard device handlers for the high-speed paper-tape reader and punch include complete interrupt subroutines for both API and PI service. Details on how the Program Interrupt Control (PIC) skip chain and the Automatic Priority Interrupt (API) channels are set up and provided in Part III of the PDP-15/10 Software System Manual. The following example of a hypothetical interrupt service subroutine is provided for general understanding of interrupt servicing.

NOTE

This example is not a complete, fully-tested interrupt service handler.

1.5.4.1 Program Interrupt Example

```

RSB
.
.
.
.PI   .LOC 0
      0           /SAVE PC, LINK, EXTEND MODE
                        /& MEM. PROT. BITS AT LOC 0.
      JMP   SKPCHN /GO TO SKIP CHAIN.
      .
SKPCHN SPFAL      /POWER FAIL FLAG TEST.
      SKP        /GO TO NEXT TEST.
      JMP* INT6  /GO TO POWER FAIL SUBRTE.
      RSF        /PAPER-TAPE READER DONE?
      SKP        /GO TO NEXT TEST.
      JMP* INT2  /GO TO PTR INTERRUPT.
      PSF        /PAPER-TAPE PUNCH DONE?
      SKP        /GO TO NEXT TEST.
      JMP* INT3  /GO TO PTP INTERRUPT.
      .
      .
      .

```

/INT6, INT2, AND INT3 ARE PART OF A TABLE
/OF INTERRUPT SERVICE ROUTINE STARTING ADDRESSES.
/AN EXAMPLE OF INT2 FOLLOWS:

INT2	PTRPIC		/15-BIT ADDRESS OF PAPER /TAPE READER SERVICE /ROUTINE.
	.		
	.		
	.		/OTHER I/O SERVICE ROUTINE /POINTERS
	.		
PTRPIC	DAC	PTRAC	/SAVE AC.
	LAC*	(0	/SAVE PC, LINK, BANK MODE /AND USER MODE IN PTROUT.
	.		
	.		
	.		/REST OF INTERRUPT HANDLED.

1.5.4.2 API Example

	RSB		/SELECT READER IN /BINARY MODE
	.		
	.		
	.		/REST OF INTERRUPT /HANDLED.
	.LOC	50	
	JMS	PTRINT	/PAPER TAPE READER /API ENTRY LOCATION
	.		
	.		
PTRINT	0		
	DAC	PTRAC	/API ENTRY. SAVE AC.
	LAC	PTRINT	/SAVE PC, LINK, BANK /MODE & USER MODE BITS.
	DAC	PTROUT	/

)

)

)

)

)

)

)

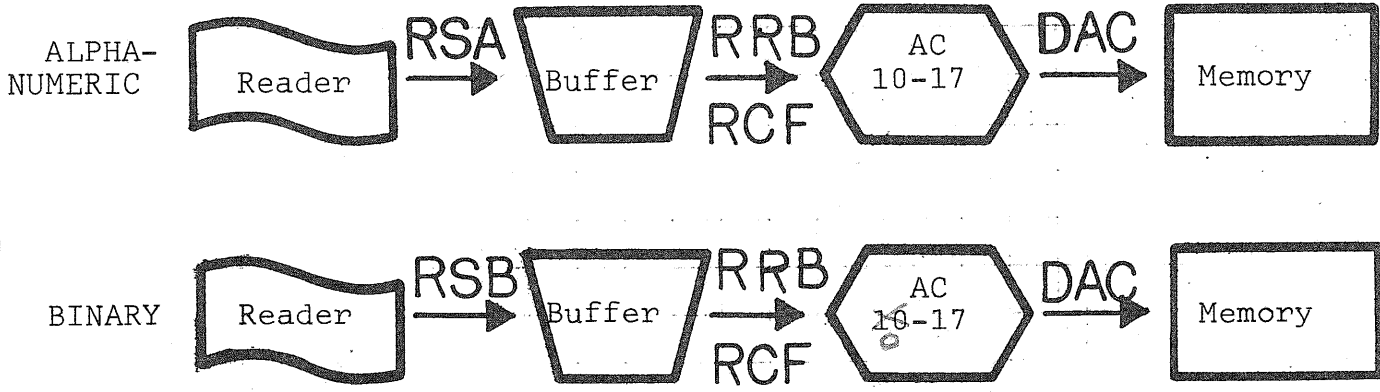
)

)

PAPER TAPE IOTS

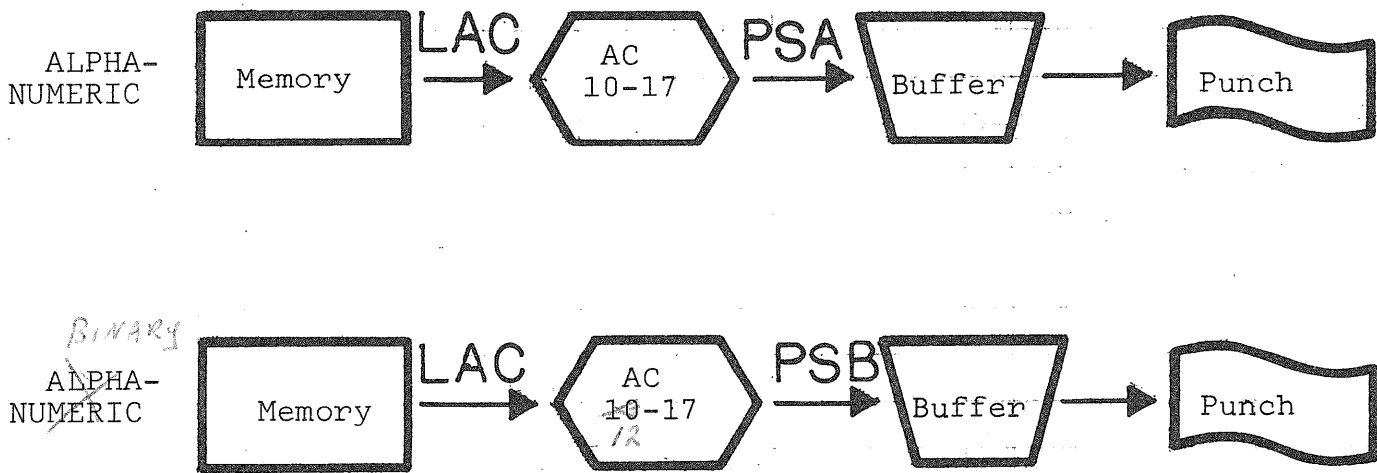
READER

INPUT



PUNCH

OUTPUT



Channel 8=1
Channel 7=0

PAPER TAPE FORMATS

Assembled Programs May Appear In Two Formats

Hardware Readin Format

1. Loaded using Hardware control logic
2. Output from assembler when use .FULL
3. Load address supplied by DATA SWITCHES on CONSOLE
4. All words loaded sequentially - no way to change where loaded
5. Read continues until a frame with CHANNEL 7 punched is detected.

Absolute Binary Format

1. Loaded using Absolute Loader
2. Output from assembler when use .ABS
3. Load address for each block supplied as first word in data block
4. All words loaded sequentially until new load address supplied by data block
5. Read continues until START BLOCK.

PAPER TAPE FORMATS

Programming tapes are supplied in one of two formats:

1. HRI - hardware read-in mode (.FULL assembly parameter)
2. BINARY OR ABS - (.ABS assembly parameter)

HRI tapes consist of 18 bit data and instructions punched in binary mode (PSB), which are loaded in sequential memory locations via the Hardware Read-In feature. The last word is an instruction which is to be executed when read (i.e. HLT or JMP). The last word is indicated by channel #7 being punched in the last frame of that word.

The load address is supplied by the address switch register.

ABS or Binary paper tapes consists of 3 basic parts:

1. ABS Loader Program punched in Hardware Read-In Format
2. Data Blocks (there may be more than one)
3. Start Block (there is only one)

The ABS Loader (Absolute Loader) is a program in HRI format. When read via the Console "Read-In" Key, it is loaded and started automatically. While executing, the Absolute Loader reads and loads the remainder of the tape. The Absolute Loader expects the tape it is reading to have a particular format containing Data and Start Blocks.

DATA BLOCK - Consists of 3 control words (Data Block Header) followed by the data to be loaded:

1. Load Address
2. Word count (not exceeding 25 and stored as a 2's complement negative number)
3. Checksum

DATA

·
·
·

START BLOCK - A two word block at the end of the tape. It is distinguished from a Data Block because bit 0 of the first word is a one (i.e. channel #6 in the first of 3 frames is punched).

1. Starting address (777777 means "HLT" rather than "JMP" to some location)
2. Dummy word (not used)

00100

00100

00101

00102

00103

00104

00105

00106

00107

00110

00111

00112

740040

750004

340112

040111

740040

750004

540111

750040

750041

000000

000001

000100

SIZE=00113

BEG

TEMP

ONE

.FULL

.LOC 100

HLT

LAS

TAD ONE

DAC TEMP

HLT

LAS

SAD TEMP

CLA!HLT

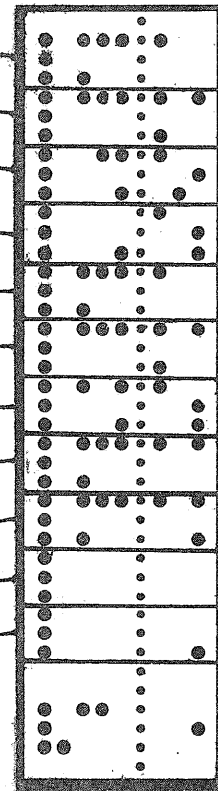
CMA!CLA!HLT

0

1

.END BEG

NO ERROR LINES

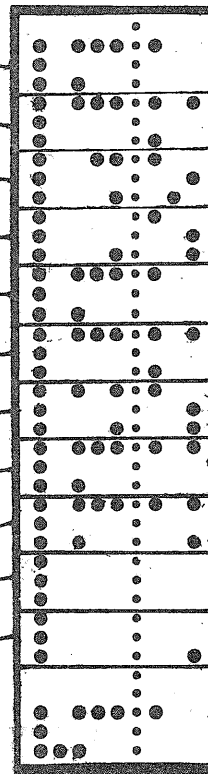


```

00100
00100 740040   BEG
00101 750004
00102 340112
00103 040111
00104 740040
00105 750004
00106 540111
00107 750040
00110 750041
00111 000000   TEMP
00112 000001   ONE
000000
      SIZE=00113
  
```

```

.FULL
.LOC 100
HLT
LAS
TAD ONE
DAC TEMP
HLT
LAS
SAD TEMP
CLA!HLT
CMA!CLA!HLT
0
1
.END
NO ERROR LINES
  
```



```

00100
00100 740040 BEG
00101 750004
00102 340111
00103 040112
00104 740040
00105 750004
00106 540112
00107 750040
00110 750041
00111 000001 ONE
000000
      SIZE=00113

```

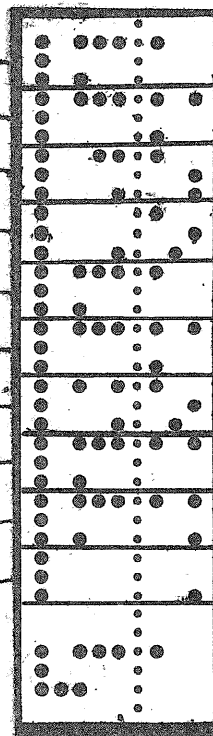
.FULL

.LOC 100

```

HLT
LAS
TAB ONE
DAC TEMP#
HLT
LAS
SAD TEMP
CLA!HLT
CMA!CLA!HLT
1

```



```

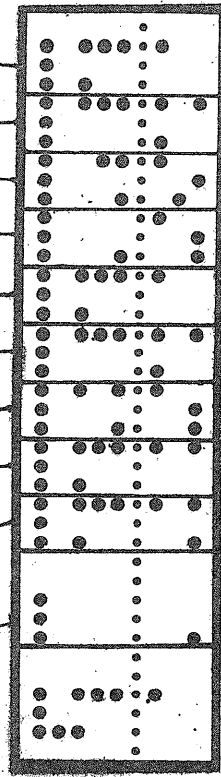
.END
NO ERROR LINES

```

```

00100 .FULL
00100 .LOC 100
00100 740040 BEG HLT
00101 750004 LAS
00102 340112 TAB (I
00103 040111 DAC TEMP#
00104 740040 HLT
00105 750004 LAS
00106 540111 SAB TEMP
00107 750040 CLA!HLT
00110 750041 CMA!CLA!HLT
00112 000000 .END
000001 *L
SIZE=00113 NO ERROR LINES

```

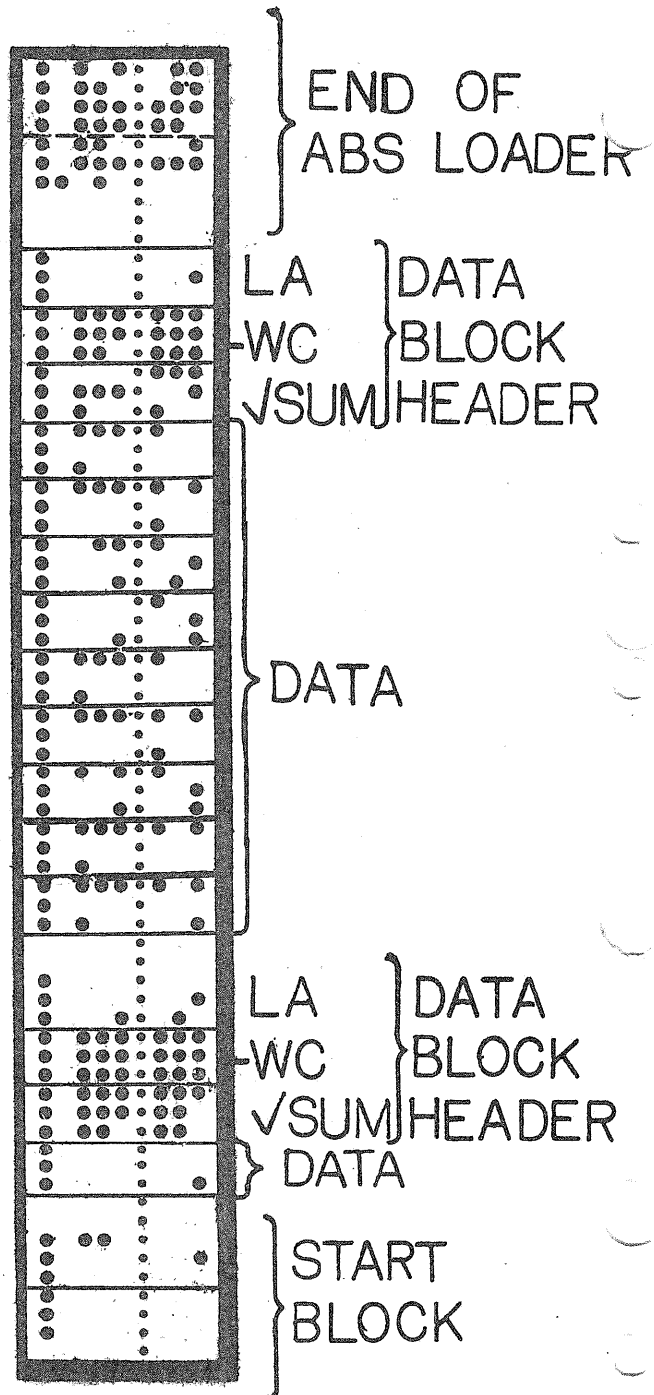


DO YOU SEE WHAT THE PROBLEM IS WHEN THIS IS LOADED?

CHECK LOCATIONS


```

00100 .ABS
00100 .LOC 100
00100 740040 BEG HLT
00101 750004 LAS
00102 340112 TAB (1
00103 040111 DAC TEMP#
00104 740040 HLT
00105 750004 LAS
00106 540111 SAB TEMP
00107 750040 CLAIHLT
00110 750041 CMAICLAIHLT
000100 .END BEG
00112 000001 *L
        SIZE=00113
        NO ERROR LINES
    
```



DOS-15 COOKBOOK

Formatting Dectapes

The format generator tape is read in under HRI mode

1. Set the address switches to 17720 (some tapes may call for 17700).
2. Depress STOP, RESET, READIN.

Conversation program begins. Place Dectape on a unit other than unit 0. Don't take up more than 2 wraps. Follow directions.

Note: Before you can use the tape, you must go to PIP to clear out the directory.

```
$PIP
DOSPIP VXX
>L,TT ← DT1
N,DT2
```

Loading DOS Into a Cold Machine

1. Mount DOS restore tape #1 on Dectape transport
Set rotary switch to 1
Set remote switch
Set write lock
2. Load DOS save-restore paper tape thru high-speed reader
Place tape in reader
Set address switches to 17720
Depress STOP, RESET and then READIN

The paper tape is read in and a conversational mode program begins.

```
INPUT: DT,
UNIT: 1,
OUTPUT: DK,
```

The program will tell you when to mount the second tape.
At the completion of loading, DOS has been placed on the disk.

3. To bring the resident Monitor into core:
Set address switches to X7637 where X = 1 for 8K
3 for 16K
5 for 24K
7 for 32K

Place the DOS bootstrap tape in the reader (usually on the same tape as the DOS SAVE & STORE paper tape)
Depress STOP, RESET, then READIN.

DOS announces itself.

➔ Note: The BOOTSTRAP stays in high core unless you cleverly manage to destroy it.

To get the monitor:

- (a) Control C (echoed as ↑C)
- (b) Set address switches to X7646 where X is same value as before: STOP, RESET, START
- (c) Set address switches to X7637 and read in bootstrap again
- (d) The whole DOS restore procedure again
- (e) Call a maintenance person

PIP

To call PIP \$PIP
 DOSPIP VXX
 >

To list the directory from the disk on the teletype:

>L TT ← DK ↵

To list the directory from the Dectape on unit 2 on the line printer:

>L_uLP ← DT2 ↵

To transfer a file from the disk to dectape:

>T_uDT1 ← DK_uTESTFL SRC ↵

>T_uDT2_uFILEA BIN ← DK_uFILL_uBIN ↵

MACRO

To call MACRO \$MACRO
 MACRO VXX
 >BN ← FILENM ↵
 OR ALTMODE

To load a program generated under .ABS or .ABSP:

1. Place tape in reader
2. Set address switches to 17720
3. Depress STOP, RESET then READIN.

This causes the absolute loader to be loaded. It starts automatically and in turn reads in the rest of the tape--that is, your program.

DOS-15 V3A000
\$L

EXAMPLE OF WRITING,ASSEMBLING AND LOADING A RELOCATABLE PROGRAM

\$EDIT

EDITOR V3A000
>OPEN SAMPLE
FILE SAMPLE SRC NOT FOUND.
INPUT
LAC (707070
HLT
.END

EDIT
>CLOSE

EDITOR V3A000
>OPEN SAMPLE
EDIT
>N
>C //START/
START LAC (707070
>L .END
.END
>A START
>P
.END START
>EXIT

DOS-15 V3A000
A LP -12

SA DK -13

SK ON

\$MACRO

BMACRO-15 V3A000
>BL←SAMPLE
END OF PASS 1
SIZE=00003 NO ERROR LINES
BMACRO-15 V3A000
>↑C

	PAGE	1	SAMPLE SRC	
00000 R	200002	R	START	LAC (707070
00001 R	740040	A		HLT
	000000	R		.END START
00002 R	707070	A *L		
	SIZE=00003			NO ERROR LINES

DOS-15 V3A000
LOAD

BLOADER V3A000
>P←SAMPLE
P SAMPLE SRC 77634
↑S↑S

REAL TIME CLOCK

References: Volume 1 Processor Handbook 6-42
System Reference Manual B-14

The Real Time Clock option provides a user with time reference capability for accounting purposes, periodic interrupts and interval timing. The clock produces clock pulses at the rate of:

or a) 60 times a second (every 16.7 ms) for 60 Hz systems
b) 50 times a second (every 20 ms) for 50 Hz systems
(the standard clock works off the line frequency--other clocks are available to produce clock pulses at user defined rates).

When the clock is enabled (CLON), every clock pulse generates a request for a break at the completion of the current instruction. When the break is granted by the CPU, the content of memory location 000007 is incremented by 1. Location 000007 is the clock counter register. As long as the clock is enabled, the process of location 7 being incremented at each clock tick continues.

When the content of location 7 overflows (i.e. is incremented from 777777 to 000000), the clock flag is set to 1. This condition may be checked for by the use of a Skip IOT (CLSF). Note also that the clock flag is interfaced to the PI and API systems so that if interrupts are enabled when the clock flag is set, an interrupt request will be made.

Three IOT instructions are associated with the clock:

CLON	700004	-Clock On -Enable the clock...increment location 000007 every clock tick -Clear the clock's flag
CLSF	700001	-Skip on Clock Flag Set...the next instruction is skipped if the clock's flag is set
CLOF	700044	-Clock Off -Disable the clock...do not increment location 000007 -Clear the clock's flag.

Since the clock counter register is memory location 000007, its contents may be modified by a program. A standard technique for using the clock is to preset the contents of location 7 to the complement of the desired time count (in ticks) and then to enable the clock (and the interrupt system if interrupts are to be used). The clock flag will be raised (and an interrupt occur) at the end of the specified time period. For instance, to raise the flag after:

1 second, set location 000007 to 777704 ($-60_{10} = -74_8$)

5 seconds, set location 000007 to 777324 ($-300_{10} = -454_8$).

Notice that it is the 2's complement that is used. This example and the following ones assume 60 ticks per second for the clock.

REAL TIME CLOCK

To check for an interval of 1 second by checking the clock's flag, the following sequence can be used:

```

LAW -74          /74 ticks = 1 second
DAC 7           /or DAC* (7 if program is not in page 0
                /                    or bank 0

CLON            /enable clock--start incrementing
CLSF           /check for clock overflow
JMP .-1        /not yet
next instruction /get here after 1 second

```

To check for an interval of 1 second via interrupts under the PI system, the following sequence may be used:

```

.LOC 0
0
JMP* .+1
TIMERT

/Main Routine
.LOC 10200
:
LAW -74          /set clock for
DAC* (7         /6010 ticks
CLON            /enable clock,clear flag
ION             /enable PI interrupts
:
:               continue with 1 second's worth of program
:

.LOC 20500

TIMERT DAC ACSAVE /save AC as it was at time of interrupt
:
:               process clock interrupt--this may involve resetting
:               location 7 or even disabling
:               the clock
:

LAC* (0         /pick up the return address from location 0
DAC RETURN#
:
:               other instructions are necessary here so that when we
:               leave TIMERT the system looks as it did before the
:               interrupt occurred.
:
JMP* RETURN     /go back to where we were interrupted

```

REAL TIME CLOCK

To check for an interval of 1 second via interrupts under the API system, the following sequence may be used:

```
.LOC 2
TIMER                /address of clock routine

.LOC 51
JMS* 2

.LOC 10200
:
:
LAW -74              /set clock counter to -60 ticks
DAC* (7
CLON                 /enable clock,clear clock flag
LAC (400000          /enable interrupts
ISA                  /from the API system
:
:                   continue with 1 second's worth of program
:

.LOC 20500
TIMER 0
DAC ACSAVE#          /save AC as it was at time of interrupt
:
:                   process clock interrupt
:
:                   other instructions are necessary here so that when we
:                   return to the interrupted routine, the system looks as
:                   it did before the interrupt occurred.
:
JMP* TIMER           /go back to where we were interrupted
```


REAL TIME CLOCK

NOTES

1. The clock continues to count up from zero after overflow occurs. At overflow detection, however, the clock counter is usually reinitialized or the clock is disabled.
2. To enable the clock:
Use the CLON instruction (make sure the console clock switch is OFF--front down; if the console is locked, then CLON will enable the clock no matter what position the console clock switch is in).

3. To disable the clock:
or a) Use the CLOF instruction
b) Turn the console CLOCK switch ON (rear half depressed; this will have an effect only if the console is not locked).
4. Depressing the RESET switch on the console clears the clock's flag and disables the clock.

.TITLE PROGRAM TO MOVE AC LIGHTS CENTER,OUT,BACK

```

/
/
/
10100 A                    .LOC 10100
10100 A 700002 A          START    IOF                    /TURN PI OFF
10101 A 754000 A                    CLAI/CLL
10102 A 705504 A                    ISA                    /TURN API OFF (ISA WITH AC000)
/
10103 A 210143 A          SETUP    LAC (002000            /SET BIT 7 BECAUSE ROTATE RIGHT
10104 A 050141 A                    DAC LHM                /BEFORE DISPLAYING
10105 A 210144 A                    LAC (000200            /SET BIT 10 BECAUSE ROTATE LEFT
10106 A 050142 A                    DAC RHM                /BEFORE DISPLAYING

```

/MANUFACTURE THE DISPLAYED ACCUMULATOR VALUE
/ BY TAKING EACH HALF AND ROTATING IT IN THE APPROPRIATE
/DIRECTION. THEN "XOR" THE HALVES FOR THE FULL VALUE.

```

/
10107 A 210141 A          FORM        LAC LH
10110 A 740020 A          MOVE/      RAR
10111 A 050141 A                    DAC LH
10112 A 210142 A                    LAC RH
10113 A 740010 A          MOVER      RAL
10114 A 050142 A                    DAC RH
10115 A 250141 A                    XOR LH
10116 A 050140 A                    DAC DISPLY#

```

/SET UP THE REAL TIME CLOCK COUNTER--LOCATION 7

```

/
10117 A 777704 A          TIMER      LAW =74                /SET UP COUNTER FOR
10120 A 070145 A                    DAC# (7                /ONE SECOND INTERVAL

```

/
/
/
.EJECT

```

/
/
/
/TURN CLOCK ON (FIRST TIME THRU) AND CLEAR CLOCK FLAG.
/SUBSEQUENT USES OF "CLON" ARE TO CLEAR THE FLAG ONLY.
/THE CLOCK KEEPS ON COUNTING AFTER IT OVERFLOWS TO 0.
/

```

```

10121 A 210140 A LAC DISPLY
10122 A 700044 A CLON
10123 A 700001 A CLSF
10124 A 610123 A JMP .-1

```

```

/
/KEEP ROTATING UNTIL THE AC=400001 OR AC=001400.
/WHEN AC EQUALS THESE VALUES CHANGE THE DIRECTION
/OF ROTATION. THIS IS DONE BY ACTUALLY CHANGING
/THE INSTRUCTION "RAL" TO "RAR" AND VICE VERSA. IT
/IS ACCOMPLISHED USING THE "XOR" INSTRUCTION AS
/RAL=740010 AND RAR=740020. SWITCHING BACK AND
/FORTH IS JUST A MATTER OF XORING WITH 000030.
/(THERE ARE OF COURSE OTHER WAYS TO SWITCH!)
/

```

```

10125 A 550146 A SAD (400001)
10126 A 610132 A JMP CHANGE
10127 A 550147 A SAD (001400)
10130 A 741000 A SKP
10131 A 610107 A JMP FORM
10132 A 210110 A CHANGE LAC MOVEL
10133 A 250150 A XOR (000030)
10134 A 050110 A DAC MOVEL
10135 A 250150 A XOR (000030)
10136 A 050113 A DAC MOVER
10137 A 610107 A JMP FORM

```

```

700044 A CLON=700044
700001 A CLSF=700001

```

```

010100 A .END START
10143 A 002000 A *L
10144 A 000200 A *L
10145 A 000007 A *L
10146 A 400001 A *L
10147 A 001400 A *L
10150 A 000030 A *L

```

SIZE=10151 NO ERROR LINES

.TITLE NON GLOBAL SUBROUTINE CALLS

```

700401 A      TSF=700401
700406 A      TLS=700406

00000 R 700002 A      START      IOF
00001 R 705514 A      ISA+10
00002 R 700416 A      TLS+10

00003 R 200043 R      LAC (TABLE
00004 R 040042 R      DAC PTR#
00005 R 777772 A      LAW =6
00006 R 040041 R      DAC COUNT#
00007 R 100033 R      JMS CRLF

00010 R 220042 R      MORE      LAC* PTR
00011 R 100026 R      JMS PRINT
00012 R 100033 R      JMS CRLF

00013 R 440042 R      ISZ PTR
00014 R 440041 R      ISZ COUNT
00015 R 600010 R      JMP MORE
                        .EXIT

00020 R 000060 A      TABLE  60
00021 R 000061 A      61
00022 R 000062 A      62
00023 R 000063 A      63
00024 R 000064 A      64
00025 R 000065 A      65

00026 R 000000 A      PRINT    0
00027 R 700401 A      TSF
00030 R 600027 R      JMP    -1
00031 R 700406 A      TLS
00032 R 620026 R      JMP* PRINT

00033 R 000000 A      CRLF     0
00034 R 760015 A      LAW 15
00035 R 100026 R      JMS PRINT
00036 R 760012 A      LAW 12
00037 R 100026 R      JMS PRINT
00040 R 620033 R      JMP* CRLF

00000 R      .END START
00043 R 000020 R #L
SIZE=00044      NO ERROR LINES

```

.TITLE GLOBAL SUBROUTINE CALLS

700401 A TSF=700401
700406 A TLS=700406

.GLOBL CRLF,PRINT

```

00000 R 700002 A     START     IOF
00001 R 705514 A             ISA+10
00002 R 700410 A             TLS+10

00003 R 200032 R             LAC (TABLE
00004 R 040027 R             DAC PTR#
00005 R 777772 A             LAW =0
00006 R 040026 R             DAC COUNT#
00007 R 120030 E             JMS* CRLF

00010 R 220027 R     MORE     LAC* PTR
00011 R 120031 E             JMS* PRINT
00012 R 120030 E             JMS* CRLF

00013 R 440027 R             ISZ PTR
00014 R 440026 R             ISZ COUNT
00015 R 600010 R             JMP MORE
                              .EXIT

00020 R 000060 A     TABLE     60
00021 R 000061 A             61
00022 R 000062 A             62
00023 R 000063 A             63
00024 R 000064 A             64
00025 R 000065 A             65

                              .END START
00030 R 000030 E +E
00031 R 000031 E +E
00032 R 000020 R +L
                              SIZE=00033     NO ERROR LINES

```

```

PAGE 1      CRLF  SRC      SUBROUTINE TO PRINT OUT CR AND LF
                                .TITLE SUBROUTINE TO PRINT OUT CR AND LF
                                .GLOBL CRLF,PRINT
                                /
                                /
00000 R 000000 A      CRLF  0
00001 R 760215 A      LAW 215
00002 R 120006 E      JMS* PRINT
00003 R 760212 A      LAW 212
00004 R 120006 E      JMS* PRINT
00005 R 620000 R      JMP* CRLF

                                .END

00006 R 000006 E *E      NO ERROR LINES
                                SIZE=00007

```

```

PAGE 1      PRINT SRC      SUBROUTINE TO PRINT CHAR ON TT
                                .TITLE SUBROUTINE TO PRINT CHAR ON TT
                                .GLOBL PRINT
                                /
                                /
700401 A      TSP#700401
700406 A      TLS#700406
                                /
                                /
                                /CHARACTER IS EXPECTED IN THE AC
                                /
00000 R 000000 A      PRINT  0
00001 R 700401 A      TSP
00002 R 600001 R      JMP  .-1
00003 R 700406 A      TLS
00004 R 620000 R      JMP* PRINT

                                .END
                                NO ERROR LINES

000000 A
                                SIZE=00005

```

DOS-15 V3A000
\$A LP -12

\$K ON

\$MACRO

BMACRO-15 V3A000
>BLG←NOGLOB
END OF PASS 1
SIZE=00044 NO ERROR LINE
↑C

DOS-15 V3A000
\$LOAD

BLOADER V3A000
>P←NOGLOB
P NOGLOB SRC 77573
↑S↑S
0
1
2
3
4
5

DOS-15 V3A000
\$MACRO

BMACRO-15 V3A000
>BLG←GLOBAL
END OF PASS 1
SIZE=00033 NO ERROR LINES
BMACRO-15 V3A000
>BL←CRLF
END OF PASS 1
SIZE=00007 NO ERROR LINES
BMACRO-15 V3A000
>BL←PRINT
END OF PASS 1
SIZE=00005 NO ERROR LINES

DOS-15 V3A000
\$LOAD

BLOADER V3A000
>P←GLOBAL,CRLF,PRINT
P GLOBAL SRC 77604
P CRLF SRC 77575
P PRINT SRC 77570
↑S↑S
0
1
2
3
4
5

.TITLE PROGRAM TO AVERAGE DECIMAL VALUES

```

/*****
/ PROGRAM ACCEPTS DECIMAL VALUES FROM THE KEYBOARD,
/ SUMS THEM AND PRINTS OUT THE DECIMAL AVERAGE (GIVEN
/ TO THE TENTHS PLACE) ON THE TELEPRINTER.
/ USER SHOULD FOLLOW EACH VALUE WITH A COMMA; TERMINATE
/ THE LINE WITH CR. FOR EXAMPLE: 3,18,29,4,(CR)
/*****/

```

```

700301 A KSF#700301
700312 A KRB#700312
700406 A TLS#700406

```

.GLOBL CRLF,PRINT

```

00000 R 707762 A BEGIN DBA /RUN IN PAGE MODE
00001 R 700002 A IOF /TURN OFF PI
00002 R 705514 A ISA+10 /AND API INTERRUPT SYSTEMS

00003 R 700416 A INIT TLS+10 /INITIATE PRINT-GET FLAG
00004 R 140117 R DZM COUNT# /COUNT# # OF VALUES
00005 R 140120 R DZM FINAL# /FINAL# SUM OF VALUES
00006 R 120125 E JMS# CRLF /ISSUE CR AND LF
00007 R 140121 R NXT DZM NUMB# /TEMPORARY LOCATION

00010 R 700301 A NEXT KSF
00011 R 600010 R JMP .-1
00012 R 700312 A KRB
00013 R 540127 R SAD (215
00014 R 600033 R JMP ALLDUN /CR MEANS LAST VALUE
00015 R 540130 R SAD (254 /", " SEPARATES VALUES
00016 R 600030 R JMP DUN
00017 R 500131 R AND (17 /GET OCTAL NUMBER
00020 R 040124 R DAC TEMP#

00021 R 200121 R LAC NUMB
00022 R 653122 A MUL
00023 R 000012 A 12
00024 R 641002 A LACO /RESULT SMALL..IN MQ
00025 R 340124 R TAD TEMP /ADD ON LAST DIGIT
00026 R 040121 R DAC NUMB /SAVE IN NUMB
00027 R 600010 R JMP NEXT

00030 R 200121 R DUN LAC NUMB /GET THIS VALUE
00031 R 340120 R TAD FINAL /ADD IT TO THE SUM
00032 R 040120 R DAC FINAL

00033 R 440117 R ISZ COUNT /KEEP TRACK OF HOW MANY VALUES
00034 R 600007 R JMP NXT

```



```

00035 R 200117 R ALLDUN LAC COUNT /GET # OF VALUES
00036 R 040041 R DAC ,+3 /STORE FOR DIVISION
00037 R 200120 R LAC FINAL /GET SUM
00040 R 653323 A IDIV
00041 R 000000 A 0
00042 R 040123 R DAC REMAINM /REMAINDER RETURNED IN AC
00043 R 641002 A LACQ /GET QUOTIENT IN HQ
00044 R 040122 R DAC QUOT#

/
00045 R 735000 A CLX
00046 R 200122 R LAC QUOT
00047 R 653323 A NEXXT IDIV
00050 R 000012 A 12
00051 R 050112 R DAC DIGIT,X /STORE IN TABLE
00052 R 641002 A LACQ /PICK UP QUOTIENT
00053 R 741200 A SNA /CONTINUE IF NOT 0
00054 R 600057 R JMP DUNN /STOP WHEN QUOTIENT = 0
00055 R 737001 A AXR 1 /PUT IN TABLE FOR LATER OUTPUT
00056 R 600047 R JMP NEXXT

/
00057 R 120125 E DUNN JMS* CRLF /ISSUE CR AND LF

/
00060 R 210112 R DUNNN LAC DIGIT,X /TO MAKE ASCII
00061 R 340132 R TAD (260
00062 R 120126 E JMS* PRINT /GET NEXT CHARACTER--WORK
00063 R 737777 A AXR -1 /WAY BACK UP TALBE
/ARE WE DONE--WHEN XR=0, YES

00064 R 724000 A PXA
00065 R 740100 A SMA
00066 R 600060 R JMP DUNNN

/
00067 R 200133 R LAC (256 /ISSUE DECIMAL POINT "."
00070 R 120126 E JMS* PRINT

/
00071 R 200117 R LAC COUNT
00072 R 040100 R DAC REM
00073 R 200123 R LAC REMAIN
00074 R 653122 A MUL
00075 R 000012 A 12
00076 R 641002 A LACQ
00077 R 653323 A IDIV
00100 R 000000 A REM 0
00101 R 641002 A LACQ
00102 R 340132 R TAD (260
00103 R 120126 E JMS* PRINT
00104 R 120125 E JMS* CRLF
00105 R 750004 A LAS /WANT TO CONTINUE OR EXIT?
00106 R 740200 A SZA /IF SWITCHES ARE 0, THEN EXIT
00107 R 600004 R JMP INIT /NON-ZERO--CONTINUE WITH ANOTHER
.EXIT

/
00112 R A DIGIT .BLOCK 5

/
000000 R .END BEGIN
00125 R 000125 E *E
00126 R 000126 E *E
00127 R 000216 A *L
00130 R 000254 A *L
00131 R 000017 A *L
00132 R 000260 A *L
00133 R 000256 A *L

```

PROGRAM INTERRUPT FACILITY (PI)

References: Volume 1 Processor Handbook 5-11
System Reference Manual 3-8

Preface: The Program Interrupt Facility increases the efficiency of input/output operations by freeing a program from the necessity of constantly monitoring device flags. When PI is enabled and a peripheral device becomes available or completes a transfer, the PI automatically interrupts the program sequence and causes a "JMS 000000" to occur. A subroutine at location 000000 may then sense the device flags to determine which of the devices caused the interrupt, service the device, and return to the main program.

The running time of programs using input and output routines is primarily made up of the time spent waiting for an I/O device to accept or transmit information. Specifically, this time is spent in loops such as:

```
TSF          /SKIP ON FLAG  
JMP .-1
```

Waiting loops waste a large amount of computer time. In those cases where the computer can be doing something else while waiting, these loops may be removed and useful routines included to use the waiting time. This sharing of a computer between two tasks is often accomplished through the program interrupt facility, which is standard on all PDP-15 computers. The program interrupt facility allows certain external conditions to interrupt the computer program. It is used to speed the processing of I/O devices or to allow certain alarms to halt program execution and initiate another routine.

Each of the input/output devices has associated with it a device flag which is set to 1 whenever the device has completed a transfer and is ready for another. When the Program Interrupt Facility is enabled, the setting of the device flag (connected to PI) causes a program interrupt request. When PI is disabled, program interrupts do not occur, although device flags may be set.

When the interrupt is granted, PI is disabled automatically, the main instruction sequence is suspended and the hardware executes a "JMS 000000". This causes the contents of the Program Counter (the address of the next instruction that was to be executed) to be stored in location 000000 and the instruction in location 000001 to be executed.

The routine entered due to the interrupt is responsible for finding and servicing the device that caused the interrupt. Usually, the instruction in location 000001 is a JMP to a sequence of code called a SKIP CHAIN which determines which device's flag caused the interrupt and then jumps into a service routine for that specific device.

PROGRAM INTERRUPT FACILITY

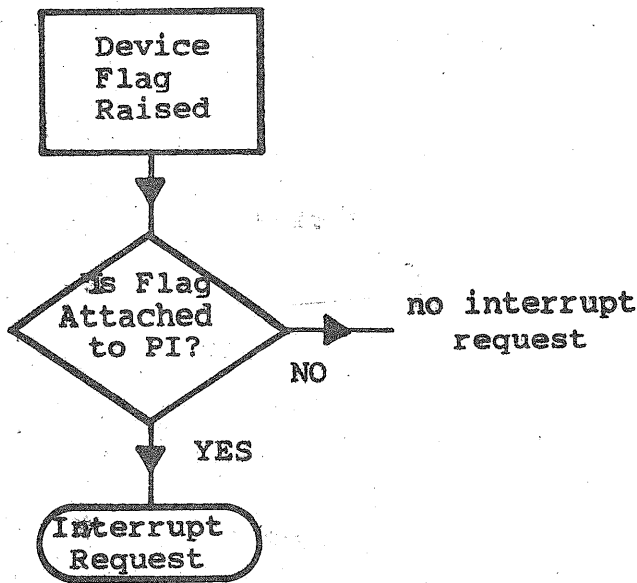
The individual service routine then handles the condition causing the interrupt, reenables the Program Interrupt system and resumes mainline program execution by JMPing to the location pointed to by location 000000.

The IOT instructions used to program the PDP-15 for program interrupts are:

ION	700042	-Interrupt ON -Enable PI interrupts
IOF	700002	-Interrupt OFF -Disable PI interrupts
<i>Vol 7 proc 6-49</i> RES	707742	-Restore...i.e. set up for the restoration of the Link, Page/Bank mode, Memory Protect bit from the pointer word given in the next indirectly referenced instruction.

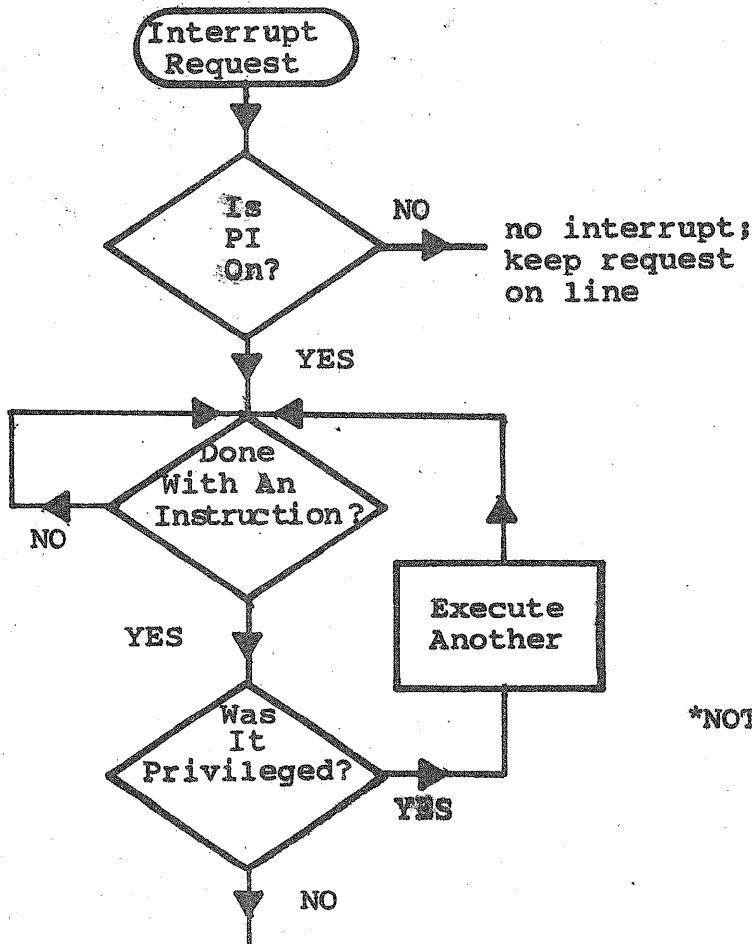
Use of the interrupt system allows a mainline routine, referred to as the BACKGROUND PROGRAM, to execute without wasting a large amount of time in waiting loops while I/O devices are assembling and transmitting information. The interrupt service routine, called a FOREGROUND PROGRAM, is entered automatically whenever an I/O device requires servicing under program control.

REQUESTING AN INTERRUPT



-- requests for program interrupts are made when flags are raised for devices tied to PI.

GRANTING AN INTERRUPT



-- when the CPU receives a request for an interrupt, it must decide if that interrupt can be granted.

PI interrupts will be granted only

- 1) if PI is on
- 2) between instructions
- 3) after a non-privileged instructions (privileged instructions: IOTs, JMS, CAL, XCT, NORM)

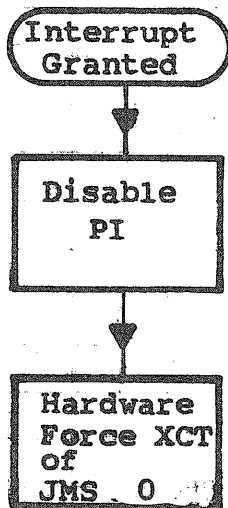
*NOTE: The following have priority over PI:

- 1) Data Channel Transfers
- 2) Clock breaks for updating location 00007
- 3) API hardware interrupts

PI Interrupt Granted*

CPU PROCESS OF AN INTERRUPT

-- an interrupt consists of:
 1) disabling PI
 2) executing a JMS 000000



Recall that information is also stored in bits 0,1, and 2 on a JMS. Location 000000

Bit 0 = Link at time of interrupt
 1 = Page/Bank Mode Indicator
 2 = Memory Protect Indicator

The word in location 000000 has the following format:

0	1	2	3	17
Link	Page Bank	Memory Protect	15 bit address of the instruction that was to be executed at time of interrupt	

PROGRAM INTERRUPT FACILITY

For example, suppose we have a PDP-15 system with high speed paper tape reader and punch, teletype and clock and that is all. The following gives a skip chain that could be used.

```

0          000000
1          JMP SKPCHN

          :
          :
MAINLINE ROUTINE
          :
          :

SKPCHN    CLSF          /did the clock cause the interrupt?
          SKP           /no
          JMP CLOCK     /yes--go to clock service routine
/
          KSF           /did the keyboard cause it?
          SKP           /no
          JMP KYBD      /yes--go to keyboard service routine
/
          TSF           /did the teleprinter?
          SKP           /no
          JMP TPRINT   /yes--go to the teleprinter service routine
/
          RSF           /did the paper tape reader?
          SKP           /no
          JMP PTREAD   /yes--go to reader service routine
/
          PSF           /did the paper tape punch?
          JMP ERROR    /no--illegal interrupt occurred-go to error routine
          JMP PUNCH    /yes--go to paper tape punch service routine

CLOCK    DAC ACSAVE
          :
          LAC ACSAVE
          ION
          RES
          JMP* 000000
/
KYBD     DAC ACSAVE
          :
          JMP* 000000
/
TPRINT   DAC ACSAVE
          :
          JMP* 000000
/
PTREAD   DAC ACSAVE
          :
          JMP* 000000
/
PUNCH    DAC ACSAVE
          :
          JMP* 000000

```

PROGRAM INTERRUPT FACILITY

An interrupt grant will cause the computer to perform the following operations automatically:

1. The PI system is disabled.
2. The contents of the PC is stored at memory location 000000.
3. The "JMP SKPCHN" in location 000001 is executed.

(note that steps 2 and 3 are the equivalent of executing "JMS 000000")

The SKPCHN routine then determines the source of the interrupt and passes control to the appropriate device handler.

The device handler (interrupt service routine) then performs the following operations:

1. The contents of the Accumulator (and any other registers which will be used) is saved.

2. The interrupt is processed --

-determine whether flag was raised due to completion of transfer or an error condition

input

-store data transferred in and clear flag

-determine if more is to be input (if so, initiate it)

output

-determine if more data is to be output (if so, initiate it)

-clear flag

3. Restore the Accumulator (and any other registers used and therefore saved).

4. Turn the interrupt system back on (if further interrupts are to be allowed).

5. Set up for the restoration of the Link, Page/Bank mode, Memory Protect mode. The "RES" instruction primes the system for this restoration, although it does not actually occur until the next indirectly referenced instruction is executed (and then it is done using the contents of bits 0,1 and 2 of the pointer word).

6. Return to the mainline program via a "JMP* 000000" instruction (recall that the updated PC was stored in location 000000).

PROGRAM INTERRUPT FACILITY

NOTES

1. Instructions like CLSF, KSF and PSF are skip-on-flag instructions. There are Skip IOTs for every device in the interrupt system. Because of the predominance of skip instructions in the instruction sequence which determines the source of an interrupt request, it is often called a SKIP CHAIN.

A skip chain may be enlarged to test for almost any number of device flags, provided that high-speed devices which retain information for a relatively short period of time are tested near the top of the skip chain, so that the chain may be traversed and the high-speed devices serviced before the information is lost. High-speed devices should never be required to wait for service while a long skip chain is traversed.

Notice that the order in which the Skip IOTs are placed in the skip chain actually determines the priority of a device. If two devices have their flags raised simultaneously, the device whose Skip IOT appears closest to the top of the skip chain will be serviced first.

2. It is possible that the SKIP CHAIN will not be in page 0 or bank 0, in which case a "JMP SKPCHN" instruction in location 1 won't allow you to get there. Instead:

```
0 000000
1 JMP* 2
2 SKPCHN
```

will allow you to get to SKPCHN because a 15 bit address is picked up from location 2.

3. Similarly, it may be that the individual device service routines will not be located in the same page or bank as the SKIP CHAIN, and therefore will have to be entered indirectly:

```
SKPCHN  KSF
        SKP
        JMP* VKB
        TSF
        SKP
        JMP* VTP
        JMP ERROR

VKB     KYBD
VTP     TPRINT
```

4. With some devices, error condition flags set to 1 will also generate interrupts. It is therefore the service routine's responsibility to determine if the interrupt was caused by the completion of a transfer or by the existence of an error condition. For example, an interrupt may be caused by the paper tape reader when:

- a) it has read a character and has assembled it in its buffer
- b) it has attempted to read the tape but finds a no tape condition.

Some error conditions may be checked using the IORS instruction. In other cases, devices have their own status registers indicating errors (e.g. MT, DK, DP, DT).

AUTOMATIC PRIORITY INTERRUPT (API)

References: Volume 1 Processor Handbook 6-46
System Reference Manual 3-10

Overview: The Automatic Priority Interrupt system option increases the capability of the PDP-15 to handle transfers of information to and from input-output devices. API identifies an interrupt device directly, without the need of a SKIP CHAIN routine for flag checking. Multi-level interrupts are permissible where a device of higher priority supersedes an interrupt already in progress. These functions increase the speed of the input-output system and simplify the programming. In this way devices (especially high speed devices) can be serviced efficiently.

The API option increases the I/O handling capabilities of the PDP-15 by adding eight levels of priority servicing (0 - 7) and associating 32 channels with these eight levels. The highest four levels of priority, i.e. 0, 1, 2, 3 are assigned to hardware devices. The lower four levels, i.e. 4, 5, 6, 7 are for software purposes.

Of the 32 API channels, 4 are assigned to the software levels 4 - 7. The remaining 28 channels are available for use by the hardware levels 0 - 3. Each of the four hardware levels may have eight devices (channels) tied to it, up to the total of 28 for the four levels. This is strictly a hardware limitation imposed by cable lengths and circuit delays, and attempts to circumvent this restriction will create needless problems.

Each of the 32 channels is assigned to a specific memory location called the Break Address. The break addresses are locations 40 - 77 in page 0, bank 0. Each device tied to API is associated with a specific channel (and therefore break address) and a specific priority level. The table below gives the standard assignments. The channel assignments should remain fixed for software compatibility, but the suggested priority level may be changed (rewiring needed) at the discretion of the user.

API ADDRESS

API Channel	Break Address	Standard Device	Suggested Priority Level
0	40	Software channel 0	4
1	41	Software channel 1	5
2	42	Software channel 2	6
3	43	Software channel 3	7
4	44	DECtape (TC15)	1
5	45	MagTape (TC59)	1
6	46		
7	47		
10	50	Paper Tape Reader (PC15)	2
11	51	Clock Overflow (KW15)	3
12	52	Power Fail (KF15)	0
13	53		0
14	54	Graphics (VT15/VP15)	2
15	55	Card Readers (CR15/CR03B)	2
16	56	Line Printer (LP15)	3
17	57	A/D (AD15/AF01)	0
20	60	DB99A/DB98A	3
21	61		
22	62	Data Phone (DP09A)	2
23	63	DECdisk (RF15)	1
24	64	Diskpack (RP15)	1
25	65	Plotter (XY15)	1
26	66		
27	67		
30	70	Scanners (DC01-ED) as needed use 70-77	3
31	71	UDC15	
32	72	ADC15	
33	73		
34	74	LT19 & LT15 Teleprinter	3
35	75	LT19 & LT15 Keyboard	3
36	76		
37	77		

Each device, when granted service via the API facility, sends its specific break address to the computer. This address, which will normally contain a JMS instruction to the device service routine, will then be executed by the computer. This type of interrupt service eliminates the need for time consuming flag search routines, and extensive core use for interrupt handling routines, by automatically determining which device requested service and providing immediate entry to the proper service routine.

Higher priority devices will be able to interrupt lower priority routines upon sending and having a request granted. The priority of devices multiplexed on the same priority level is determined by the relative position of the devices on the I/O bus. The first device on the bus having highest priority at that level, the second having second highest priority, etc.

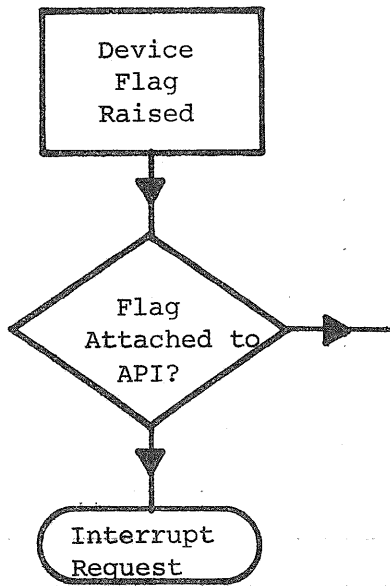
The entire API facility can be enabled or disabled by a single IOT instruction. There is no way to enable or disable specific priority levels. However, for some devices there are instructions to disconnect itself from the API facility.

In addition to the above, there are two special features in the API facility. These are:

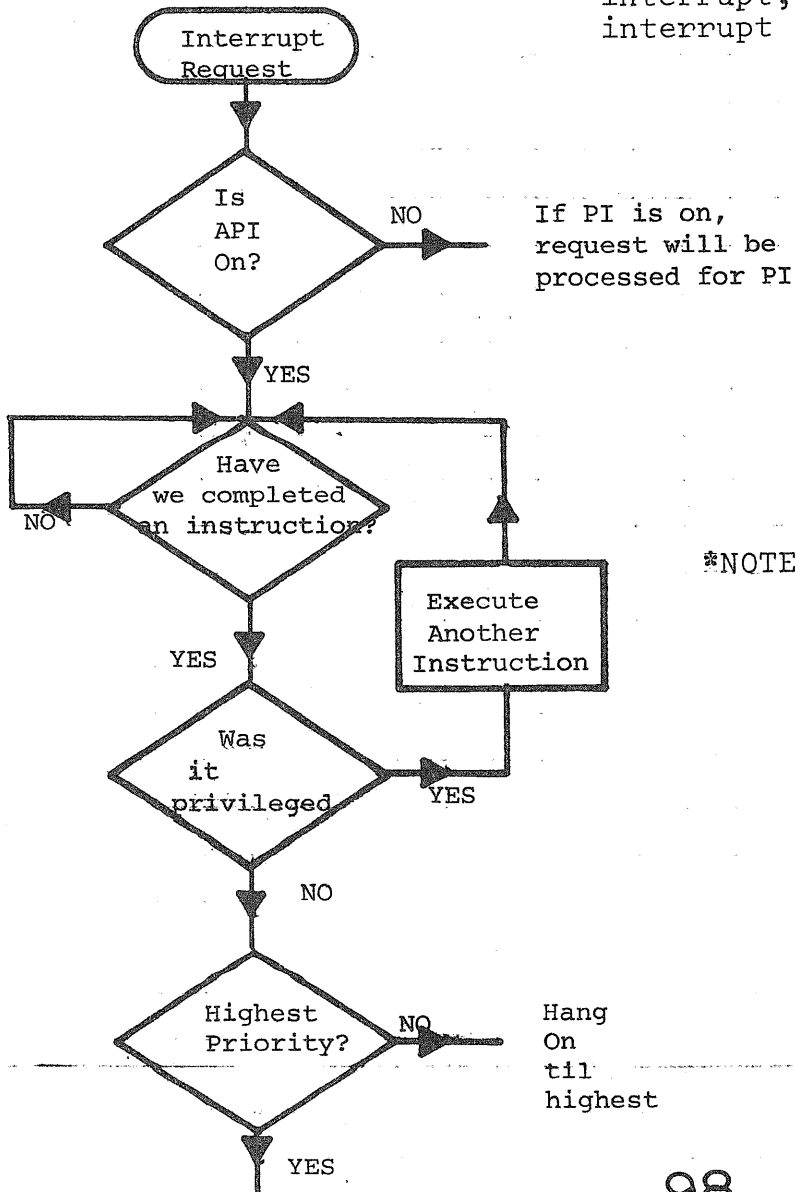
- a. The CAL Instruction - Execution of a CAL instruction with the API facility enabled automatically sets priority level 4 thereby shutting out software requests of a lower priority until this level is released.
- b. Program Interrupt - A program interrupt, from any I/O device connected to the computer, sets priority level 3. This occurs whether or not the API facility is enabled. This causes all devices on priority level 3, all software requests and program interrupts to be shut out until the level is released.

Special care must be taken in the programming of the API option to take account of these two features.

REQUESTING AN API INTERRUPT - requests for API interrupts are made when device flags, which are tied to the API system, are raised.



GRANTING AN INTERRUPT - when the CPU receives a request for an interrupt, the CPU must decide if that interrupt can be granted under API.



*NOTE: The CPU decides if it can grant an API interrupt in response to the API request. The following have priority over API:

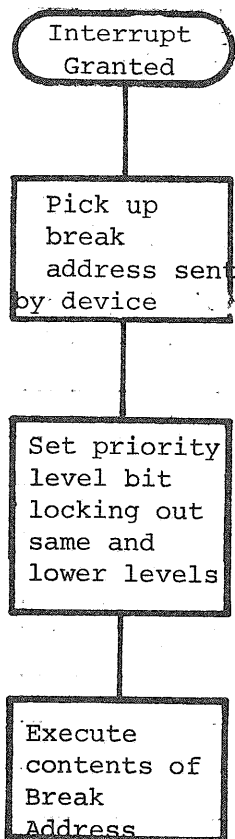
- 1) Data Channel Transfers
- 2) Clock breaks for updating location 000007.

PROCESSING AN API INTERRUPT -- in processing an API interrupt, the CPU

1) sets a priority level bit to inhibit interrupts from channels of the same or lower level of priority

2) picks up the break address sent by the device

3) does a hardware forced XCT of the contents of the break address

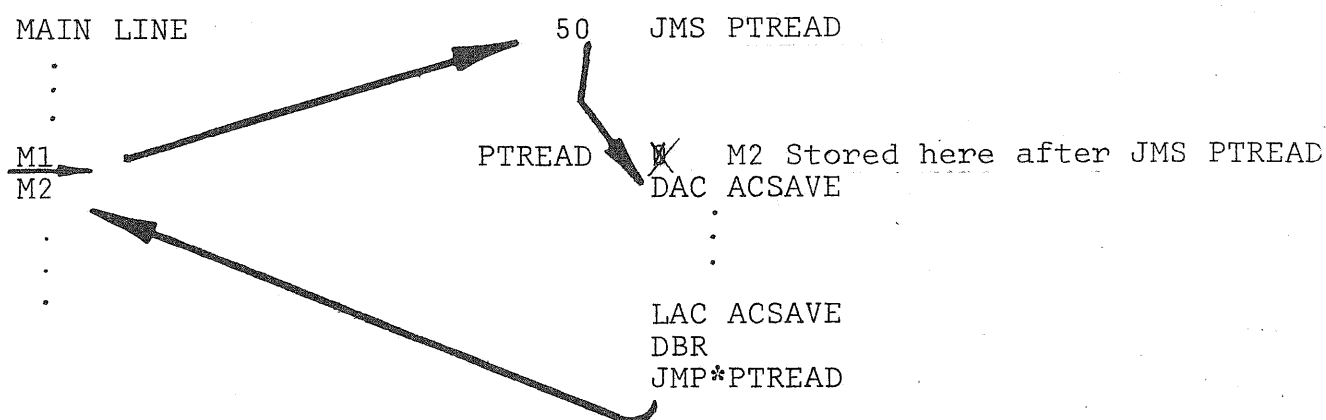


The IOT instructions used to program the API system are:

- DBK 703304 -Debreak
-Reset the highest priority level bit so that operations may be carried out on same or lower level.
- RES 707742 -Restore
-Set up to restore the link, page/bank mode and memory protect mode on next indirect instruction.
- DBR 703344 -Debreak and Restore
- ISA 705504 -Initiate Selected Activity
-Used to initiate software level interrupts and to raise the priority level of an operating program.

SINGLE DEVICE INTERRUPT PROGRAMMING

When programming a system with only one possible source of interrupts, say the paper tape reader, the handling of an interrupt is straight forward and simple.



When an API interrupt is granted to the Paper Tape Reader, it sends to the CPU its break address (50). The contents of location 50 is "XCT"ed. Since an XCT instruction does not change the PC, the PC that gets stored in location PTREAD is the updated PC from the mainline program (location M2). The interrupt is processed by PTREAD and it does a debreak to release level 2 (its priority level) and then Jumps back to the mainline program using location PTREAD as a pointer.

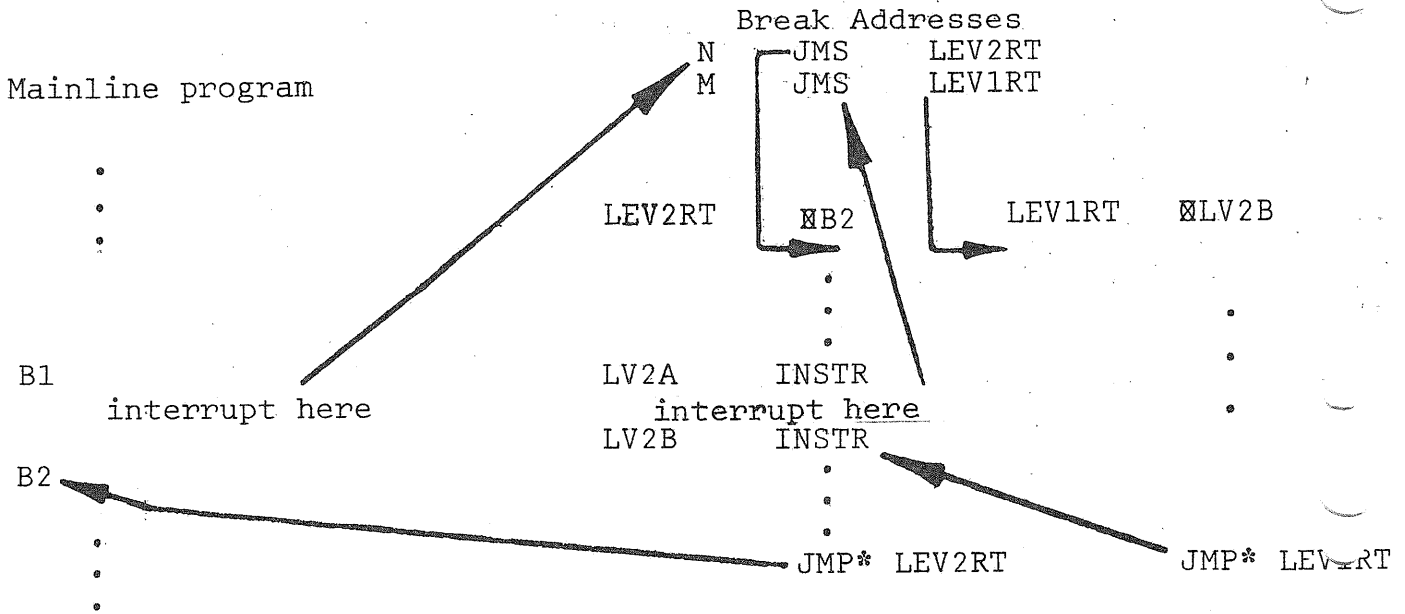
MULTIPLE DEVICE PROGRAMMING

When there is more than one device attached to API, handling an interrupt simply requires setting up the associated break address with a JMS to the interrupt service routine.

50	JMS	PTREAD
51	JMS	CLOCK
52	JMS	PWRFL
PTREAD	Ø	
	DAC	ACSAV1
	LAC	ACSAV1
	DBR	
	JMP*	PTREAD
CLOCK	Ø	
	DAC	ACSAV2
	LAC	ACSAV2
	DBR	
	JMP*	CLOCK
PWRFL	Ø	
	DAC	ACSAV3
	LAC	ACSAV3
	DBR	
	JMP*	PWRFL

With this system there is no need for polling because when a device is granted an interrupt, it sends to the CPU its associated break address, which can then contain a JMS to a routine to handle that device.

While in the service routine for one device, a higher priority interrupt may be granted. This presents no problem because of the manner in which return addresses are stored.



In this example, a level 2 device has break address N while a level 1 device has break address M. The routine to service the level 2 interrupt is LEV2RT and the routine to service the level 1 interrupt is LEV1RT.

When the mainline routine is interrupted by the level 2 device between instructions in B1 and B2, the PC is pointing to B2. The level 2 device sends its break address of N to the CPU so that the JMS LEV2RT instruction may be "XCT"ed. This causes the address B2 (the contents of the PC) to be stored in location LEV2RT and control passed to location LEV2RT + 1. If while we are at priority 2 in routine LEV2RT an interrupt is requested by a device at level 1 that interrupt can be granted because of its higher priority. Suppose it is granted between instructions at LV2A and LV2B. Then the PC contains the address LV2B. The level 1 device sends its break address of M to the CPU so that the JMS LEV1RT instruction may be executed. This causes the address LV2B to be stored at location LEV1RT and control passed to location LEV1RT + 1. The level 1 routine processes the level 1 interrupt. When it executes the JMP* LEV1RT instruction control is passed back to the level 2 routine at the point we left off, LV2B. The LEV2RT routine may now resume its operation. When the JMP* LEV2RT instruction is executed, we go back to the mainline routine at location M2 where we left off and continue from there.

RELOCATION RULES

- A. IF ADDRESS IS A NUMBER (NOT A SYMBOL) THE ADDRESS IS ABSOLUTE.
- B. IF THE ADDRESS IS A SYMBOL WHICH IS DEFINED BY A DIRECT ASSIGNMENT STATEMENT (I.E. =) AND THE RIGHT-HAND SIDE OF THE ASSIGNMENT IS A NUMBER, ALL REFERENCES TO THE SYMBOL WILL BE ABSOLUTE.
- C. IF A USER LABEL OCCURS WITHIN A BLOCK OF CODING THAT IS ABSOLUTE, THE LABEL IS ABSOLUTE.
- D. VARIABLES, UNDEFINED SYMBOLS, EXTERNAL TRANSFER VECTORS, AND LITERALS GET THE SAME RELOCATION AS WAS IN EFFECT WHEN .END WAS ENCOUNTERED IN PASS 1.
- E. IF THE LOCATION COUNTER (.LOC PSEUDO OP) REFERENCES A SYMBOL WHICH IS NOT DEFINED IN TERMS OF AN ABSOLUTE ADDRESS, THE SYMBOL IS RELOCATABLE.
- F. ALL OTHERS ARE RELOCATABLE.

PAGE 1 RULES EXA

/EXAMPLE OF RELOCATION RULES FOR LINKING LOADER

```

                                .GLOBAL SUB
000005 A      A=5
000000 R      B=START
000000 R 200005 A      START   LAC A
000001 R 200010 A      LAC 10
000002 R 200000 R      LAC B
000003 R 200000 R      LAC START
000004 R 220010 E      LAC* SUB
000005 R 200011 A      LAC (100
000006 R 200012 A      LAC (START

000005 R
000005 R 200010 A      START1  .LOC START+5
000006 R 200005 R      LAC 10
000007 R 000000 R      LAC START1
                                START

000000 A
000000 A 200000 R      .LOC 0
000001 A 200005 R      LAC START
000002 A 200002 A      LAC START1
000003 A 000000 R      START2  LAC START2
                                START

000000 A      .END
00010 A 000010 E *E
00011 A 000100 A *L
00012 A 000000 R *L

```

SIZE=00013 NO ERROR LINES

.↑C

DOS-15 V3A000
\$A LP -12

\$K ON

\$PAGE ON

\$LOAD

LOADER V3A000
>P←RELOC
P RELOC SRC 77614
↑S↑Q

DOS-15 V3A000
\$DUMP

DUMP V3A000
>77614-77637

DUMP V3A000
>↑C

DOS-15 V3A000
\$BANK ON

\$LOAD

BLOADER V3A000
>P←RELOC
P RELOC SRC 77614
↑S↑Q

DOS-15 V3A000
\$DUMP

DUMP V3A000
>77614-77637

DOS-15 V3A000
\$

```

00000 R 200013 R LAC SYMBOL
00001 R 200017 R LAC B#
00002 R 040014 R DAC SYMBOL+1
00003 R 040017 R DAC B
00004 R 220013 R LAC* SYMBOL
00005 R 160011 A DZM* 11
00006 R 220010 A LAC* 10
00007 R 200010 A LAC 10
00010 R 200021 R LAC (END
00011 R 200022 R LAC (END=SYMBOL
00012 R 040020 R DAC C
00013 R 005000 A SYMBOL 5000
00014 R 000000 A 0
00015 R 000004 A 4
00016 R 000005 A END 5
00021 R 000016 R *L .END
00022 R 000003 A *L
SIZE=00023 1 ERROR LINES
  
```

PAGE LOAD

77614-77637

77610	000000	000000	000000	000000	207627	207633	047630	047633
77620	227627	160011	220010	200010	207635	207636	047634	005000
77630	000000	000004	000005	000000	000000	077632	000003	000762

BANK LOAD

77614-77637

77610	000000	000000	000000	000000	217627	217633	057630	057633
77620	237627	160011	220010	200010	217635	217636	057634	005000
77630	000000	000004	000005	000000	000000	077632	000003	000762

)

)

)

)

)

)

)

DOS-15 V3A000
PIP

DOSPIP V3A000

>L TT COPIA BIN ← DTI

06-MAR-75
DIRECTORY LISTING
346 FREE BLKS
66 USER FILES
110 SYSTEM BLKS
COPIA BIN 455 1

>L TT COPIA BIN ← DK (P)

06-MAR-75
DIRECTORY LISTING (PES)
2202 FREE BLKS
32 USER FILES
34 USER BLKS
COPIA BIN 1567(2) 1 06-MAR-75 1567 100

>↑C

DOS-15 V3A000
\$A DTI -14

\$A TT -12

\$DUMP

DUMP V3A000
455#

455#

0	015000	646531	262601	000004	000005	000070	071033	412450
10	034150	074623	230204	400000	000000	000004	040404	000001
20	000000	000000	040404	001005	000001	000000	040404	000000
30	002004	000010	015000	317540	050404	000025	777736	000004
40	040404	000012	002005	000011	050404	000025	777736	000005
50	040302	000012	600010	000067	040710	000000	407716	022600
60	230710	000025	420564	052033	005000	520247	230710	000067
70	474741	071640	232700	000010	000010	000011	001005	776773
100	TO	367	CONTAINS	000000				
370	000000	000000	000000	000000	000000	000000	000000	777777

.IODEV 4,5

00000 R 000004 A *G
 00001 R 000001 A *G
 00002 R 000000 A *G
 00003 R 000000 A *G

.INIT 4,0,0
 CAL+0*1000 4&777
 1
 0+0
 0

00004 R 001005 A *G
 00005 R 000001 A *G
 00006 R 000000 A *G
 00007 R 000000 A *G

.INIT 5,1,0
 CAL+1*1000 5&777
 1
 0+0
 0

00010 R
 00010 R 002004 A *G
 00011 R 000010 A *G
 00012 R 000025 R *G
 00013 R 777736 A *G

START .READ 4,2,BUFF,34
 CAL+2*1000 4&777
 10
 BUFF
 .DEC
 -34

00014 R 000004 A *G
 00015 R 000012 A *G

.WAIT 4
 CAL 4&777
 12

00016 R 002005 A *G
 00017 R 000011 A *G
 00020 R 000025 R *G
 00021 R 777736 A *G

.WRITE 5,2,BUFF,34
 CAL+2*1000 5&777
 11
 BUFF
 .DEC
 -34

00022 R 000005 A *G
 00023 R 000012 A *G

.WAIT 5
 CAL 5&777
 12

00024 R 600010 R

JMP START

00025 R A

BUFF .BLOCK 42

00067 R 000000 A

ENDMRK 0

000010 R
 SIZE=00070

.END START
 NO ERROR LINES

↑C

DOS-15 V3A000
SA DK -14

\$DUMP

DUMP V3A000
>1567#

1567#

0	015500	646031	262601	000004	000005	000070	071033	412450
10	034150	074623	230204	400000	000000	000004	040404	000001
20	000000	000000	040404	001005	000001	000000	040404	000000
30	002004	000010	015500	317040	050404	000025	777736	000004
40	040404	000012	002005	000011	050404	000025	777736	000005
50	040302	000012	600010	000067	040710	000000	407716	022600
60	230710	000025	420564	052033	005500	517547	230710	000067
70	474741	071640	232700	000010	000010	000011	001005	776773
100	000001	001567	000000	000000	000000	000000	000000	000000
110	TO	367	CONTAINS 000000					
370	000000	000000	000000	000000	000000	000000	777777	777777

RADIX 50₈ VALUES

X--		-X-		--X	
A	003100	A	000050	A	000001
B	006200	B	000120	B	000002
C	011300	C	000170	C	000003
D	014400	D	000240	D	000004
E	017500	E	000310	E	000005
F	022600	F	000360	F	000006
G	025700	G	000430	G	000007
H	031000	H	000500	H	000010
I	034100	I	000550	I	000011
J	037200	J	000620	J	000012
K	042300	K	000670	K	000013
L	045400	L	000740	L	000014
M	050500	M	001010	M	000015
N	053600	N	001060	N	000016
O	056700	O	001130	O	000017
P	062000	P	001200	P	000020
Q	065100	Q	001250	Q	000021
R	070200	R	001320	R	000022
S	073300	S	001370	S	000023
T	076400	T	001440	T	000024
U	101500	U	001510	U	000025
V	104600	V	001560	V	000026
W	107700	W	001630	W	000027
X	113000	X	001700	X	000030
Y	116100	Y	001750	Y	000031
Z	121200	Z	002020	Z	000032
%	124300	%	002070	%	000033
.	127400	.	002140	.	000034
0	132500	0	002210	0	000035
1	135600	1	002260	1	000036
2	140700	2	002330	2	000037
3	144000	3	002400	3	000040
4	147100	4	002450	4	000041
5	152200	5	002520	5	000042
6	155300	6	002570	6	000043
7	160400	7	002640	7	000044
8	163500	8	002710	8	000045
9	166600	9	002760	9	000046
#	171700	#	003030	#	000047

GE

PROG1 SRC

00000 R
 000000 A
 000000 A
 SIZE=02000
 .BLOCK 2000
 .END
 NO ERROR LINES

PAGE 1 PROG2 SRC

00000 R
 000000 A
 000000 A
 SIZE=07500
 .BLOCK 7500
 .END
 NO ERROR LINES

PAGE 1 PROG3 SRC

00000 R
 000000 A
 000000 A
 SIZE=00500
 .BLOCK 500
 .END
 NO ERROR LINES

PAGE 1 PROG4 SRC

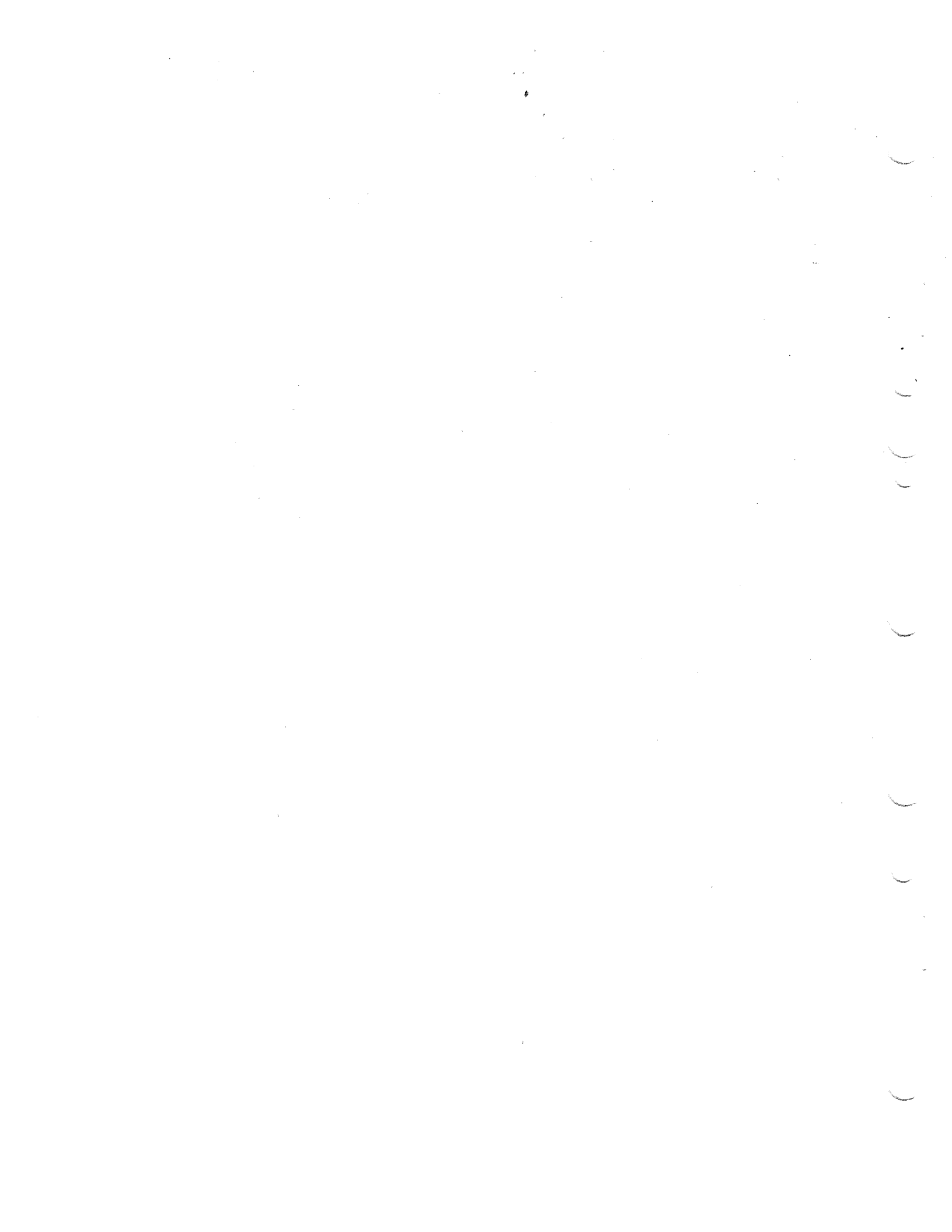
00000 R
 000000 A
 000000 A
 SIZE=07740
 .BLOCK 7740
 .END
 NO ERROR LINES

PAGE 1 PROG5 SRC

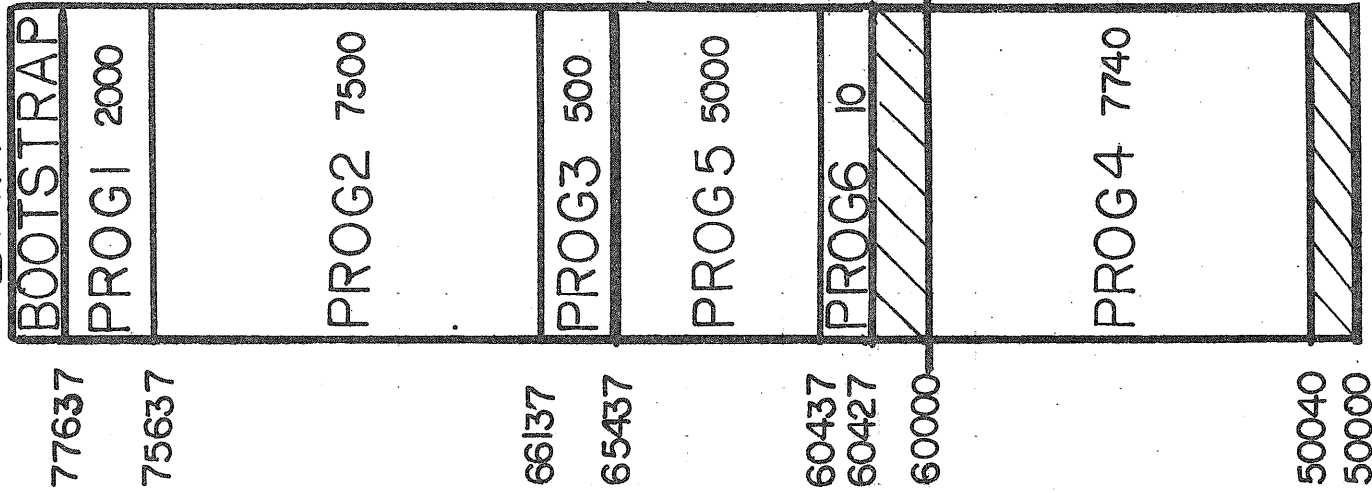
00000 R
 000000 A
 000000 A
 SIZE=05000
 .BLOCK 5000
 .END
 NO ERROR LINES

PAGE 1 PROG6 SRC

00000 R
 000000 A
 000000 A
 SIZE=00010
 .BLOCK 10
 .END
 NO ERROR LINES



BANK



\$LOAD

BLOADER V3A000
 >P-PROG1,PROG2,PROG3,PROG4,PROG5,PROG6
 P PROG1 SRC 75637
 P PROG2 SRC 66137
 P PROG3 SRC 65437
 P PROG4 SRC 50040
 P PROG5 SRC 60437
 P PROG6 SRC 60427
 ↑StC

DOS-15 V3A000

\$PAGE ON

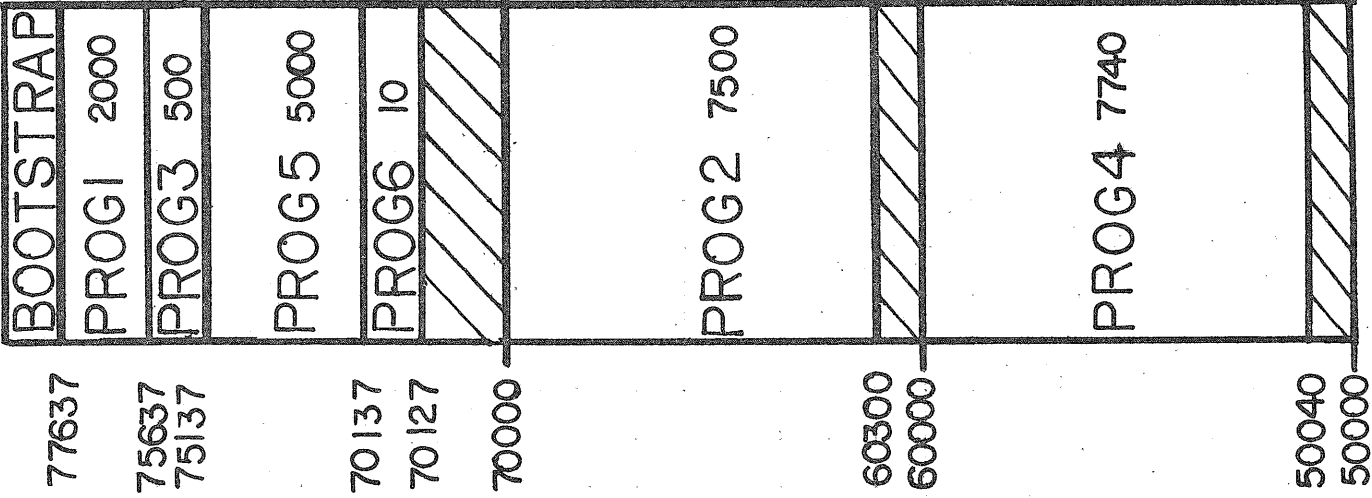
\$LOAD

LOADER V3A000
 >P-PROG1,PROG2,PROG3,PROG4,PROG5,PROG6
 P PROG1 SRC 75637
 P PROG2 SRC 60300
 P PROG3 SRC 75137
 P PROG4 SRC 50040
 P PROG5 SRC 70137
 P PROG6 SRC 70127
 ↑StC

DOS-15 V3A000

\$

PAGE



FORTRAN-IV AND MACRO

In previous chapters, MACRO calling sequences have been given for OTS and Science Library Subprograms. This general form is used in a MACRO program to call any FORTRAN external subroutine or function. A FORTRAN program may also invoke MACRO subprograms. The method for each type of linkage is given below.

➔ INVOKING MACRO SUBPROGRAMS FROM FORTRAN

A FORTRAN program may invoke any MACRO program whose name is declared in a MACRO .GLOBL statement. The MACRO subprogram must also include the same number of open registers as there are arguments. These will serve as transfer vectors for arguments supplied in the FORTRAN CALL statement or function reference. A FORTRAN-IV program and the MACRO subprogram it invokes are shown below.

FORTRAN	MACRO
C TEST MACRO SUBR	.TITLE MIN
C READ A NUMBER(A)	.GLOBL MIN, .DA
1 READ(1,100)A	MIN 0 / entry/exit
100 FORMAT(E12.4)	JMS* .DA / general get
C NEGATE THE NUMBER	/ argument
C AND PUT IT IN B	/ (OTS)
CALL MIN(A,B)	JMP .+2+1 / jump around
C WRITE OUT NUMBER(B)	argument
WRITE(2,100)B	registers
STOP	MIN1 .DSA 0 / ARG1
END	MIN2 .DSA 0 / ARG2
	LAC* MIN1 / first word of A
	DAC* MIN2 / store at B
	ISZ MIN1 / point to second word
	ISZ MIN2 / of A and B
	LAC* MIN1 / second word of A
	RAL / sign bit = 1
	CML /
	RAR /
	DAC* MIN2 / store in second
	JMP* MIN word of B
	.END exit

The FORTRAN statement CALL MIN(A,B) is expanded by the compiler to:

```
00013 JMS* MIN           / to MACRO subprog
00014 JMP $00014
00015 .DSA A
00016 .DSA B
$00014 = 00017
```

When the FORTRAN-IV program is loaded, the addresses (plus relocation factor) of A and B are stored in registers 15 and 16, respectively. When the MACRO program invokes .DA, these addresses are stored in MIN1 and MIN2 and the values themselves are accessed by indirect reference.

Arguments are, as described above, transmitted by .DA using a single word. Bits 3-17 contain the 15-bit address of the first word. Bits 0-2 serve as flag. FORTRAN uses bit 0 to indicate that the word specifying the argument contains the address of a word containing the address of the first word of the argument. The MACRO argument word always contains the address of the first word of the argument. For array name arguments (unsubscripted), the address of the fifth word of the array descriptor block is given with bit 0 on.

For external functions, the MACRO subprogram must return with a value in the AC (LOGICAL, INTEGER), AC-MQ (DOUBLE INTEGER) or in the floating accumulator (REAL or DOUBLE PRECISION).

➡ INVOKING FORTRAN SUBPROGRAMS FROM MACRO

The MACRO calling conventions for FORTRAN subprograms are: the name of the subprogram must be declared as global; there must be a jump around the argument address; and the number and mode of arguments in the call must agree with those of the subprogram. This form is shown below.

```
.TITLE  MACPRG
.GLOBL  SUBR
JMS*    SUBR
JMP     .+N+1           / jump around arguments ignored by .DA
.DSA    ARG1           / address of first argument - bit 0 set to 1
.DSA    ARG2           / indicates indirect reference
.
.
.DSA    ARGN
```

When the subprogram is compiled, a call is generated to .DA which performs the transmission of arguments from MACRO. The beginning of a subroutine might be expanded as follows.

C	TITLE SUBR
000000	SUBROUTINE SUBR(A,B)
000001	CAL 0
000002	JMS* .DA
000003	JMP \$000002
000004	.DSA A
\$ 000002 = 000005	.DSA B

If a value is to be returned by the subroutine, it is most convenient to have this be one of the calling arguments. An external function is called in the same manner as a subroutine but returns a value in the AC (single integers), AC-MQ (double integers), or floating accumulator (real and double-precision). To store the AC, the MACRO program uses a DAC instruction. Values from the floating accumulator may be stored via the OTS routines .AH (real) and .AP (double-precision). For FPP systems, values are returned in a hardware accumulator and stored with an FST instruction.

➔ COMMON BLOCKS

FORTRAN COMMON blocks (and block-data subprograms) may be linked to MACRO programs. When the MACRO program is loaded, global symbols are first sought in the user and system libraries. Any remaining are matched, where possible, to COMMON block names. This cannot be done if programs are loaded via CHAIN and EXECUTE. For example:

FORTRAN	MACRO
INTEGER A,B,C	.GLOBL NAME, .XX / .XX is name given to blank COMMON
COMMON/NAME/C	/by the FORTRAN Compiler
COMMON A,B	DZM* .XX / CLEAR A - NOTE INDIRECT REFERENCE
.	ISZ .XX / BUMP COUNTER
.	DZM* .XX / CLEAR B
.	DZM* NAME / CLEAR C

Note that if the values are REAL (two words) or DOUBLE PRECISION (three words), the MACRO program must account for the number of words when accessing specific variables.

DOS-15 and RSX-PLUS MACRO programs may also use the .CBD pseudo-op. For instance

```
BASE1 .CBD NAME 1
```

will provide the base address of the common block NAME in the word that is created and labeled BASE1; the size of the common block is 1. For blank common, use for example:

```
BASE2 .CBD .XX 2
```


)

)

)

)

)

)

)

)

.TITLE MACRO ROUTINE CALLING FORTRAN PROGRAM

.GLOBL FORT

```

/
/*****
/ PROGRAM REQUESTS USER TO INPUT THE # OF ELEMENTS TO
/ BE SUMMED IN ARRAY A(1 TO 10 ELEMENTS MAY BE SUMMED).
/ THE USER INPUTS A VALUE THAT IS ONE LESS THAN THE
/ ACTUAL NUMBER DESIRED. THE INPUT RANGE IS 0-9.
/ MACCAL CALLS FORT TO ADD 1 TO THE NUMBER TYPED.
/ FORT IN TURN CALLS MACEX TO DO THE SUMMING.
/
/*****

```

```

000000 A
000001 A

```

```

IN=0
OUT=1

```

00000 R

```

START .INIT =2,IN,RST
      .INIT =3,OUT,RST

```

00010 R

```

RST .WRITE =3,2,MES1,34
    .WAIT =3

```

00016 R

```

CHECK .READ =2,3,BUFF,3      /READS 1 CHARACTER
      .WAIT =2

```

```

00024 R 200146 R
00025 R 340151 R
00026 R 040146 R
00027 R 741100 A
00030 R 600046 R
00031 R 340152 R
00032 R 740100 A
00033 R 600046 R
00034 R 200146 R
00035 R 040055 R

```

```

CK1 LAC BUFF+2
    TAD (-60
    DAC BUFF+2
CK2 SPA
    JMP ERRMES
    TAD (-12
CK3 SMA
    JMP ERRMES
CK4 LAC BUFF+2
    DAC COUNT

```

```

/X>60, OR X=60
/X<60 SO ENTRY NOT 0-9
/X=60-12<0 MEANS X<72
/X<72 MEANS ENTRY <OR= 9
/X>OR=72 MEANS NOT 0-9

```

```

00036 R 120150 E
00037 R 600041 R
00040 R 000055 R

```

```

JMS+ FORT
JMP .+2
.DSA COUNT

```

```

00041 R 750004 A
00042 R 740200 A
00043 R 600010 R

```

```

LAS
SZA
JMP RST

```

.EXIT

00046 R

```

ERRMES .WRITE =3,2,MESERR,34
      .WAIT =3
      JMP CHECK

```

00054 R 600016 R

00055 R 000000 A

```

COUNT 0

```

```

00056 R 021000 A
00057 R 000000 A
00060 R 522632 A
00061 R 042500 A
00062 R 446344 A
00063 R 047634 A
00064 R 425011 A
00065 R 442646 A
00066 R 515012 A
00067 R 444202 A
00070 R 471012 A
00071 R 444212 A
00072 R 202352 A
00073 R 546604 A
00074 R 426444 A
00075 R 047614 A
00076 R 202131 A
00077 R 442632 A
00100 R 426352 A
00101 R 451500 A
00102 R 546372 A
00103 R 520256 A
00104 R 476531 A
00105 R 442100 A
00106 R 462231 A
00107 R 342500 A
00110 R 516531 A
00111 R 546612 A
00112 R 421005 A
00113 R 030132 A
00114 R 345225 A
00115 R 000000 A
00116 R 064000 A
00117 R 000000 A
00120 R 012000 A
00121 R 000000 A
00122 R 532031 A
00123 R 452612 A
00124 R 202232 A
00125 R 320236 A
00126 R 526504 A
00127 R 047614 A
00130 R 202450 A
00131 R 147216 A
00132 R 425344 A
00133 R 000000 A
00134 R 522632 A
00135 R 042500 A
00136 R 235405 A
00137 R 534516 A
00140 R 202371 A

```

MES1

MESERR=MES1/2*1000

0

.ASCII /TYPE IN ONE LESS THAN THE NUMBER OF ELEM

.ASCII /YOU WOULD LIKE SUMMED (0-9)./

.ASCII <15>

MESERR

BUFF=MESERR/2*1000

0

.ASCII /VALUE IS OUT OF RANGE. /

.ASCII /TYPE '0-9' ONLY/<15>

00141 R 646262 A
 00142 R 064000 A
 00143 R 000000 A

00144 R 002000 A
 00145 R 000000 A
 00146 R A

!
 BUFF ENDLOC=BUFF/2*1000
 0
 .BLOCK 2

000150 R
 000000 R

!
 ENDLOC=.
 .END START

00150 R 000150 E *E
 00151 R 777720 A *L
 00152 R 777766 A *L

SIZE=00153

NO ERROR LINES

```

001      C      FORTRAN CALLING MACRO EXAMPLE
002      C
003      C      THIS ROUTINE IS CALLED BY THE MACRO ROUTINE "MACCAL".
004      C      "MACCAL" PASSES ONE ARGUMENT, N , TO FORT
005      C      FORT ADDS 1 TO IT AND THEN...
006      C
007      C      THIS FORTRAN ROUTINE CALLS MACRO PROGRAM SUM
008      C      TO PERFORM SUMMATION OF ARRAY A.
009      C      THEN THE SUM IS PRINTED OUT AFTER TOT.
010      C
011      C      SUBROUTINE FORT(N)
012      C
013      C      INTEGER A(10),TOT
014      C      A(1)=5
015      C      A(2)=6
016      C      A(3)=7
017      C      A(4)=8
018      C      A(5)=9
019      C      A(6)=10
020      C      A(7)=11
021      C      A(8)=12
022      C      A(9)=13
023      C      A(10)=14
024      C
025      C      N=N+1
026      C      CALL SUM(A,N,TOT)
027      C      WRITE(6,1) TOT
028      C
029      1  FORMAT(1H , 'TOT=', I10)
030      C      RETURN
031      C      END

```

```

001      C      FORTRAN CALLING MACRO EXAMPLE
002      C
003      C      THIS ROUTINE IS CALLED BY THE MACRO ROUTINE "MACCAL".
004      C      "MACCAL" PASSES ONE ARGUMENT, N , TO FORT
005      C      FORT ADDS 1 TO IT AND THEN...
006      C
007      C      THIS FORTRAN ROUTINE CALLS MACRO PROGRAM SUM
008      C      TO PERFORM SUMMATION OF ARRAY A.
009      C      THEN THE SUM IS PRINTED OUT AFTER TOT*.
010      C
011      C      SUBROUTINE FORT(N)
00000 .DSA FORT
00001 JMS* .DA
00002 JMP $00002
00003 .DSA N
$00002 * 00004
012      C
013      C      INTEGER A(10),TOT
014      C      A(1)=5
00004 CMAICLA
00005 TAD (000001
00006 TAD A
00007 DAC $00007
00010 LAC (000005
$00007 * 00153
00011 DAC* XIA
015      C      A(2)=6
00012 CMAICLA
00013 TAD (000002
00014 TAD A
00015 DAC $00015
00016 LAC (000006
$00015 * 00153
00017 DAC* XIA
016      C      A(3)=7
00020 CMAICLA
00021 TAD (000003
00022 TAD A
00023 DAC $00023
00024 LAC (000007
$00023 * 00153
00025 DAC* XIA
017      C      A(4)=8
00026 CMAICLA
00027 TAD (000004
00030 TAD A
00031 DAC $00031
00032 LAC (000010
$00031 * 00153
00033 DAC* XIA
018      C      A(5)=9
00034 CMAICLA
00035 TAD (000005
00036 TAD A
00037 DAC $00037
00040 LAC (000011
$00037 * 00153
00041 DAC* XIA
019      C      A(6)=10
00042 CMAICLA

```

```

00043 TAD (000006
00044 TAD A
00045 DAC 800045
00046 LAC (000012
800045 = 00153
00047 DAC* XIA
020 A(7)=11
00050 CMAICLA
00051 TAD (000007
00052 TAD A
00053 DAC 800053
00054 LAC (000013
800053 = 00153
00055 DAC* XIA
021 A(8)=12
00056 CMAICLA
00057 TAD (000010
00060 TAD A
00061 DAC 800061
00062 LAC (000014
800061 = 00153
00063 DAC* XIA
022 A(9)=13
00064 CMAICLA
00065 TAD (000011
00066 TAD A
00067 DAC 800067
00070 LAC (000015
800067 = 00153
00071 DAC* XIA
023 A(10)=14
00072 CMAICLA
00073 TAD (000012
00074 TAD A
00075 DAC 800075
024 C
00076 LAC (000016
800076 = 00153
00077 DAC* XIA
025 NON+1
00100 LAC* N
00101 TAD (000001
00102 DAC* N
026 CALL SUM(A,N,TOT)
00103 JMS* SUM
00104 JMP 00110
00105 .DSA 400000 +A
00106 .DSA 400000 +N
00107 .DSA TOT
027 WRITE(6,1) TOT
00110 JMS* .FW
00111 .DSA (000006
00112 .DSA .1
028 C
00113 .DSA 777777
00114 JMS* .FE
00115 .DSA TOT
00116 JMS* .FF
029 1 FORMAT(1H , 'TOT=', I10)
00117 JMP 800117

```

```

00120 .DSA 241431
00121 .DSA 020130
00122 .DSA 235511
00123 .DSA 752172
00124 .DSA 235311
00125 .DSA 130540
00126 .DSA 245004
00127 .DSA 020100
000117 * 00130
#30 RETURN
00130 JMP .EX
#31 END
00131 JMP* 00000
00132 .DSA .DA
00133 .BLK 000012
00146 .DSA 000000
00146 .DSA 000012
00147 .DSA 000000
00150 .DSA 000000
00151 .DSA A
00152 .BLK 000001
00153 .BLK 000001
00154 .DSA SUM
00155 .DSA .FW
00156 .DSA .FE
00157 .DSA .FF
00160 .DSA 000001
00161 .DSA 000005
00162 .DSA 000002
00163 .DSA 000006
00164 .DSA 000003
00165 .DSA 000007
00166 .DSA 000004
00167 .DSA 000010
00170 .DSA 000011
00171 .DSA 000012
00172 .DSA 000013
00173 .DSA 000014
00174 .DSA 000015
00175 .DSA 000016
FORT 17777
* .DA 00132
N 00003
.EX 00131
A 00133
TOT 00152
XIA 00153
* SUM 00154
.i 00117
* .FW 00155
* .FE 00156
* .FF 00157

```


UPDATE

DOS-15 V3A000
\$R UPDATE

.DAT	DEVICE	UIC	USE
-15	DKA	PES	OUTPUT
-14	DKA	PES	INPUT
-12	TTA	PES	LISTING
-10	TTA	PES	SECONDARY INPUT

CREATING A LIBRARY WITH UPDATE

\$A LP -12/DK -10

\$K ON

\$UPDATE

UPDATE V3A000

>NL←USERLB

>I CRLF

>CLOSE

} USERLB CONTAINS ONE BINARY PROGRAM

UPDATE V3A000

>NL←.LIBR5

>I CRLF

>I PRINT

>CLOSE

} .LIBR5 CONTAINS TWO BINARY PROGRAMS

UPDATE V3A000

>NL←MAIN

>I AVGLOB

>I CRLF

>I PRINT

>CLOSE

} MAIN CONTAINS THREE BINARY PROGRAMS

UPDATE V3A000

>↑C

DOS-15 V3A000

\$

)

)

)

)

)

)

)

)

)

↑C

DOS-15 V3A000
\$PAGE ON

\$R LOAD

.DAT	DEVICE	UIC	USE
-5	NON	PES	USER LIBR
-4	DKA	PES	USER PROG(S)
-1	DKA	SYS	SYS LIBR

MAKING USE OF LIBRARIES
WITH THE LINKING LOADER

\$A DK -5

\$K ON

\$LOAD

LOADER V3A000

>P←AVGLOB,CRLF,PRINT
P AVGLOB SRC 77503
P CRLF SRC 77474
P PRINT SRC 77467

→ ALL THREE FILES SPECIFIED.
NONE PULLED FROM A LIBRARY.

↑S↑S
3,4,6,
4.3

1,2,3,
2.0

DOS-15 V3A000
\$LOAD

LOADER V3A000

>P←AVGLOB
P AVGLOB SRC 77503
P CRLF SRC 77474
P PRINT SRC 77467

→ ONLY ONE FILE SPECIFIED.
THE OTHER TWO PROGRAMS--CRLF AND PRINT--
ARE PULLED FROM THE USER LIBRARY .LIBR5
AUTOMATICALLY.

↑S↑S
3,4,5,
4.0

DOS-15 V3A000
\$LOAD

LOADER V3A000

>P←MAIN
P AVGLOB SRC 77503
P CRLF SRC 77474
P PRINT SRC 77467

→ THE FILE SPECIFIED IS A LIBRARY.
EVERY PROGRAM IN THE LIBRARY IS LOADED.

↑S↑S
3,4,5,
4.0

DOS-15 V3A000
\$LOAD

LOADER V3A000

>P←AVGLOB,PRINT
P AVGLOB SRC 77503
P PRINT SRC 77476
P CRLF SRC 77467

→ TWO FILES (PROGRAMS) ARE SPECIFIED.
THE THIRD PROGRAM --CRLF--
IS AUTOMATICALLY PULLED FROM THE
USER LIBRARY .LIBR5.

↑S↑S
3,4,5,

DOS-15 V3A000
\$LOAD

LOADER V3A000

>P-AVGLOB,USERLB,PRINT → THREE FILES ARE SPECIFIED.
P AVGLOB SRC 77503 THE FILE AVGLOB CONTAINS ONE PROGRAM(AVGLOB)
P CRLF SRC 77474 THE USER LIBRARY USERLB CONTAINS ONE
P PRINT SRC 77467 PROGRAM--CRLF--.
↑S↑S THE FILE PRINT CONTAINS ONE PROGRAM (PRINT).
3,4,5,
4.0

DOS-15 V3A000
\$

BATCH (NON-BOSS)

References: DOS USERS GUIDE 8-30
DOS KEYBOARD COMMAND GUIDE 11

DOS-15 V3A000
\$L

PREPARING A BATCH STREAM

\$EDIT

EDITOR V3A000
>OPEN BATSTR
FILE BATSTR SRC NOT FOUND.
INPUT
\$JOB
PIP
L LP ← DK <SCR >
L LP FILBLK BIN ← DK (P)
\$JOB
A LP -12
MACRO
BLG←ME
\$JOB
GLOAD
P←ME
\$EXIT

EDIT
>EXIT

DOS-15 V3A000
\$L

GETTING THE BATCH STREAM OUT ON PAPER TAPE

\$PIP

DOSPIP V3A000

>I PP ← DK BATSTR SRC

DOS-15 V3A000

\$L

RUNNING THE BATCH STREAM...PLACE THE PAPER TAPE IN THE READER

\$BATCH PR

DOS-15 V3A000

\$\$JOB

PIP

DOSPIP V3A000

>L LP ← DK <SCR>

>L LP FILBLK BIN ← DK (P)

>\$JOB

DOS-15 V3A000

\$A LP -12

\$MACRO

MACRO-15 V3A000

>BLG←ME

END OF PASS 1

SIZE=00025 NO ERROR LINES

MACRO-15 V3A000

>\$JOB

DOS-15 V3A000

\$GLOAD

LOADER V3A000

>P←ME

P ME SRC 77612

P LPA.15 049 77050

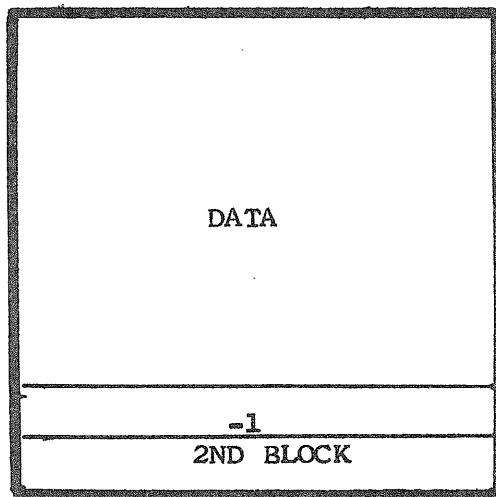
DOS-15 V3A000

\$\$EXIT

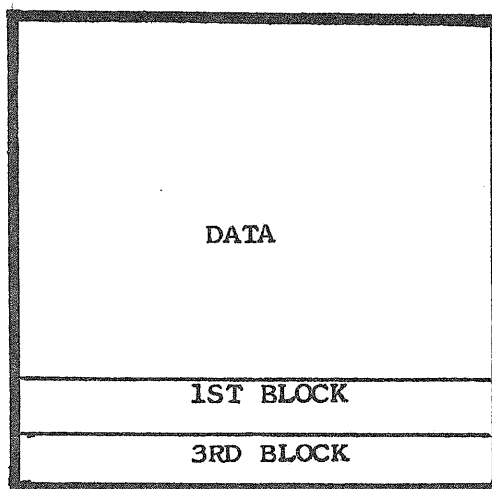
DOS-15 V3A000

\$

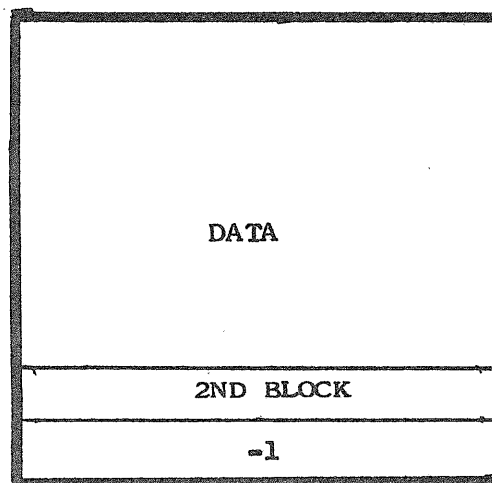
DISK FILE FORMAT FOR DOS:



1ST BLK



2ND BLK



3RD BLK

.TITLE PROGRAM TO LIST ON TT BLOCKS IN A FILE ON

/
 /
 /*****
 /

/PROGRAM LISTS ON THE TELETYPE THE BLOCKS USED BY A FILE
 /STORED ON THE DEVICE ASSOCIATED WITH DAT 5.
 /USER MUST KNOW THE STARTING BLOCK NUMBER (GIVEN BY PIP
 /ON DT DIRECTORY OR OBTAINED BY USING THE (P) SWITCH
 /FOR DIRECTORY ON DISK. [L TT @ DK (P)]
 /WHEN PROGRAM HALTS, PLACE THE NUMBER OF THE FIRST BLOCK
 /USED BY THE FILE IN THE DATA SWITCHES AND THEN HIT THE
 /CONTINUE SWITCH.

/
 /*****
 /

.IODEV 5

00000 R 707762 A

/
 START DBA /MAKE SURE IN PAGE MODP
 .INIT =3,1,0 /INITIALIZE DEVICES
 .INIT 5,0,0

/
 /*****
 /GET BLOCK NUMBER AND FILL IT IN APPROPRIATE WORD OF
 /.TRAN EXPANSION.
 /*****

00011 R 740040 A
 00012 R 750004 A
 00013 R 040017 R

HLT /READ IN BLOCK NUMBER
 LAS /FROM DATA SWITCHES
 DAC TRAN+2 /PUT IN .TRAN EXPANSION
 /BLOCK # POSITION

00014 R 100033 R

JMS PRINT /PRINT OUT BLOCK #

00015 R

TRAN .TRAN 5,0,0,BUFF,256 /DO .TRAN
 .WAIT 5

/
 /*****
 /CHECK THE LAST WORD IN THE BLOCK JUST READ IN.
 /IF THIS WORD = 777777, THEN THE BLOCK JUST READ IS THE
 /IN THE FILE.
 /IF THIS WORD NOT = 777777, THEN IT IS THE NUMBER OF TH
 /BLOCK USED BY THIS FILE.
 /*****

00024 R 200472 R
 00025 R 340473 R
 00026 R 000055 R
 00027 R 100033 R

LAC BUFF+377 /PICK UP WORD 377 IN BLK
 SAD (777777) /LAST FILE BLK(777777)?
 JMP ENDFIL /YES
 JMS PRINT /NO-PRINT OUT BLOCK #

```

/*****
/GET HERE WHEN THERE ARE STILL MORE BLOCKS TO THE FILE.
/PICK UP THE NEXT BLOCK NUMBER FROM WORD 377 OF FILE.
/PLACE IT IN THE APPROPRIATE WORD IN THE .TRAN EXPANSION
/THEN GO DO THE .TRAN
/*****

```

```

00030 R 200472 R
00031 R 040017 R
00032 R 000015 R

```

```

LAC BUFF+377 /SET UP TO TRAN IN NEXT
DAC TRAN+2
JMP TRAN /DO NEXT TRAN

```

```

/*****
/CONVERT 6 OCTAL DIGITS TO 6 ASCII CHARACTERS
/*****

```

```

00033 R 000000 A
00034 R 052000 A
00035 R 735000 A
00036 R 200474 R
00037 R 722000 A
00040 R 750000 A
00041 R 640603 A
00042 R 340475 R
00043 R 050003 R
00044 R 725001 A
00045 R 000040 R

```

```

PRINT 0
LMG
CLX
LAC (0
PAL
NXT CLA
LLS 3
TAD (00
DAC BUFF0+2,X
AXS 1
JMP NXT

```

```

/*****
/OUTPUT BLOCK #
/*****

```

```

00054 R 020033 R

```

```

.WRITE =3,3,BUFF0,0
.WAIT =3
JMP* PRINT

```

```

/*****
/ON END OF FILE, CLOSE DAT =3 AND RETURN TO MONITOR
/*****

```

```

00055 R

```

```

ENDFIL .CLOSE =3
.EXIT

```

```

00061 R 005003 A
00062 R 000000 A
00063 R A
00071 R 000015 A
00072 R 000012 A

```

```

BUFF0 005003
0
.BLOCK 6
15
12

```

```

00073 R A

```

```

BUFF .BLOCK 400

```

```

000800 R
00473 R 777777 A *L
00474 R 000000 A *L
00475 R 000000 A *L

```

```

.END START

```

```

SIZE=00476 NO ERROR LINES

```

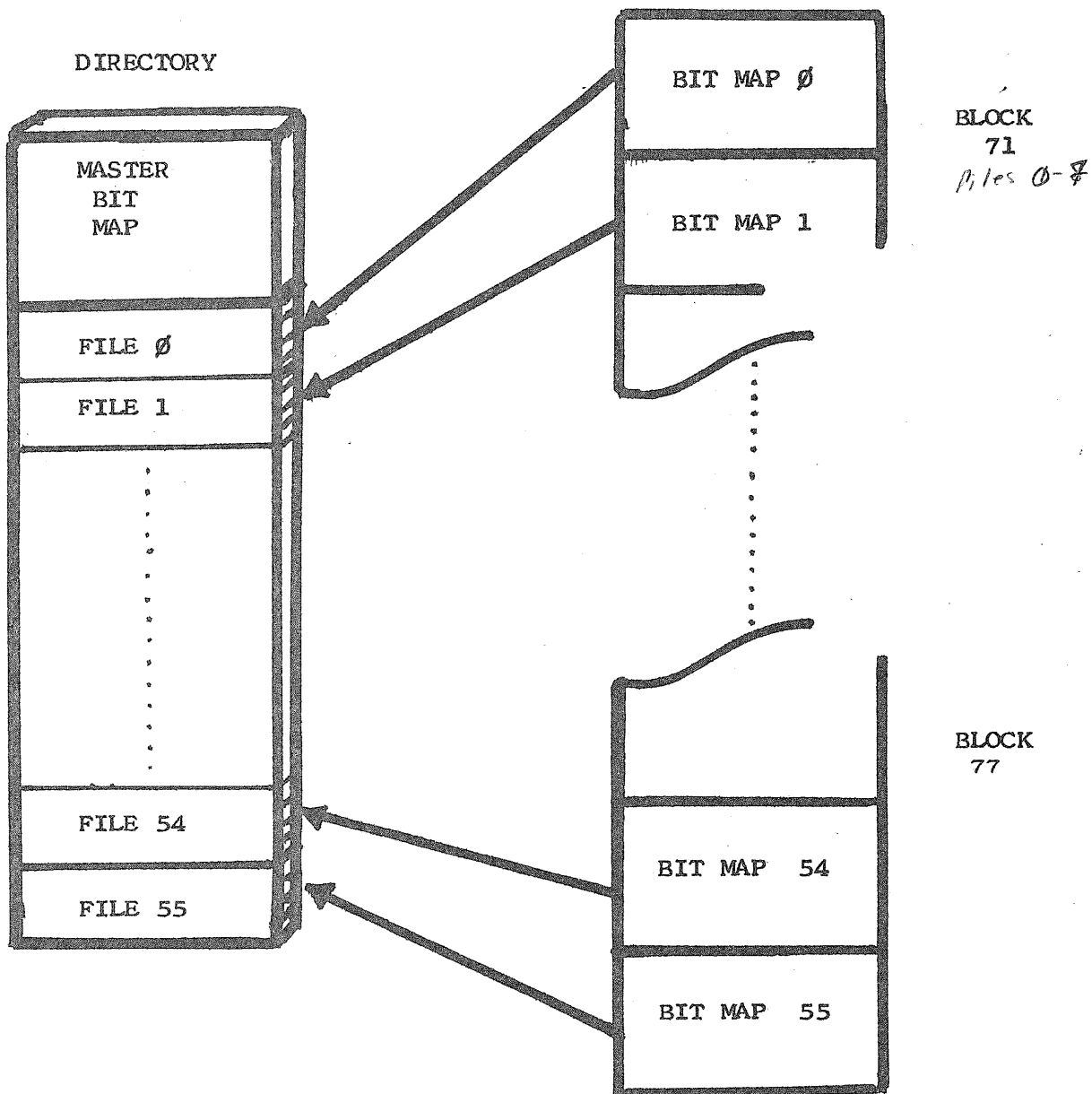
EXAMPLE OF ONE BLOCK OF A FILE:

HEADER WORD 0
HEADER WORD 1
"DATA"
HEADER WORD 0
HEADER WORD 1
"DATA"
HEADER WORD 0
HEADER WORD 1
"DATA"
EMPTY
NOT USED (padding)
NEXT BLOCK OR (-1)

BIT MAPS

BLOCKS 71 THROUGH 77 MAINTAIN INDIVIDUAL BIT MAPS FOR ALL FILES ON DECTAPE.

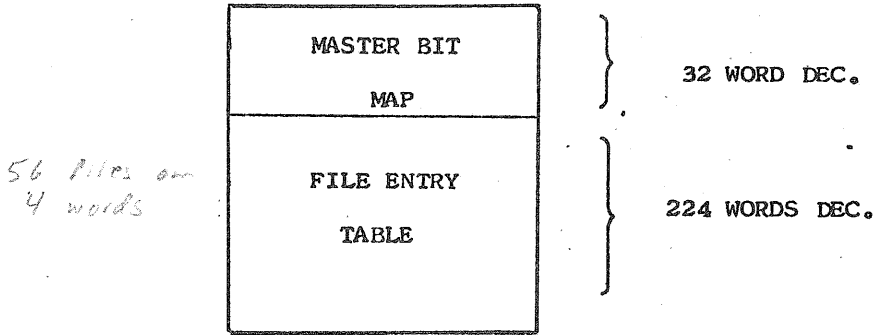
BIT MAP FILE RELATION:



THE INCLUSIVE "OR" OF ALL INDIVIDUAL BIT MAPS IS EQUIVALENT TO THE MASTER BIT MAP.

DECTAPE FILE STRUCTURE

DIRECTORY (BLK #100):



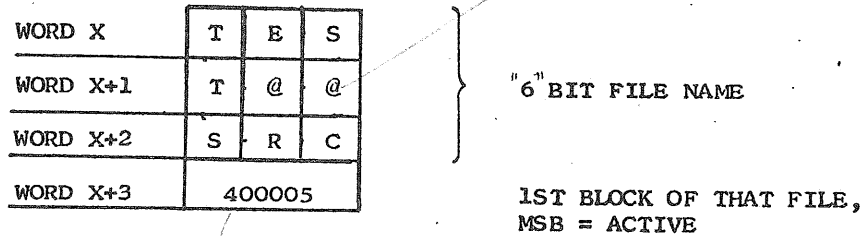
THE MASTER BIT MAP MAINTAINS A COMPLETE BLOCK BY BLOCK RECORD OF A DECTAPE BY RELATING AN INDIVIDUAL BIT TO AN OCTAL BLOCK NUMBER: E. G.

BIT POSITIONS	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17
WORD 0	0	1	2	3	4	5	6	7	10	11	12	13	14	15	16	17	20	21
WORD 1	22	23	24	25	26	27	30	31	32	33	34	35	36	37	40	41	42	43
WORD 3	44	45	.	.	.	ETC.												

file number
1 = block 0
0 = block 1000

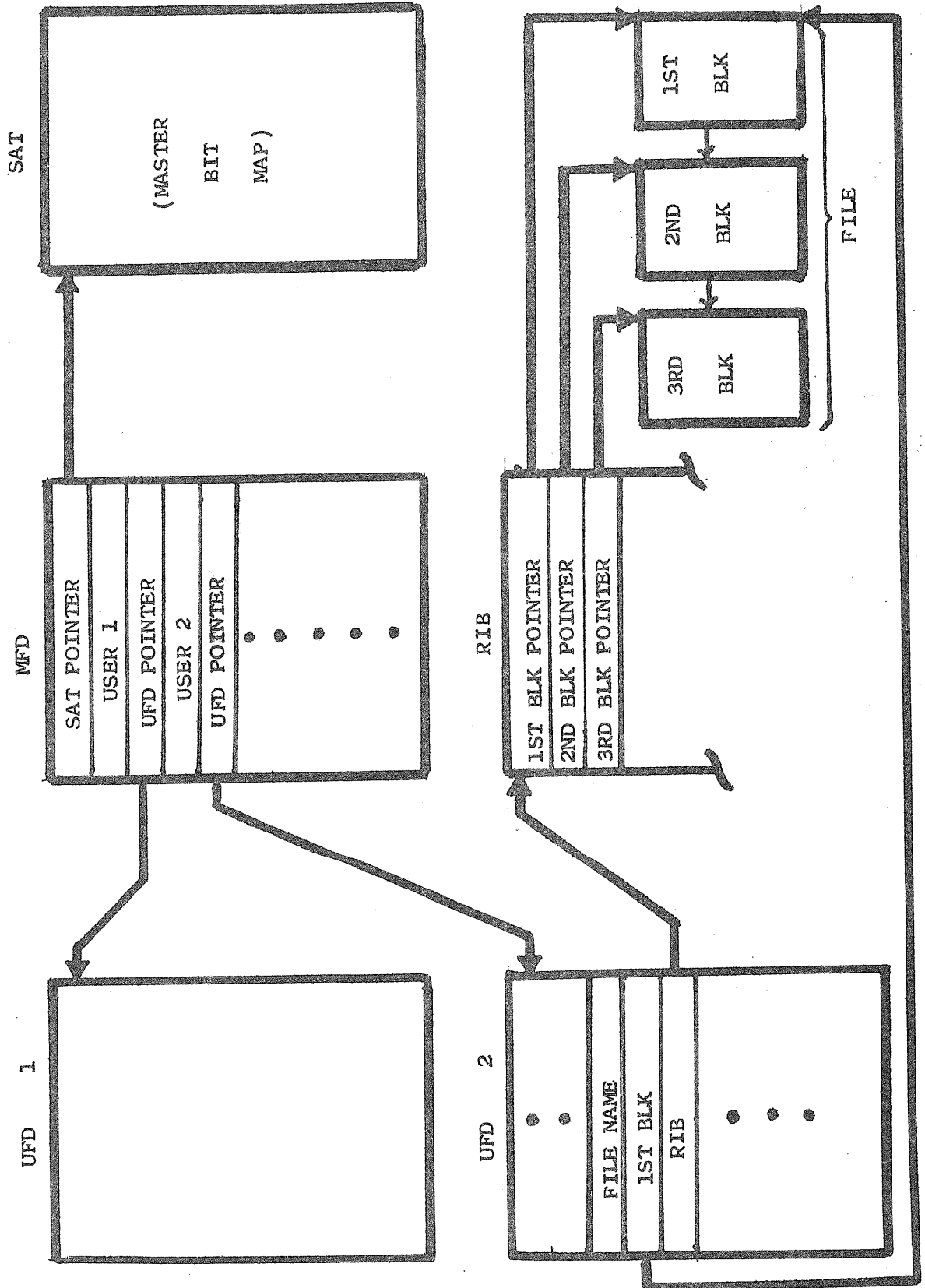
THE FILE ENTRY TABLE IS DIVIDED IN 4 WORD SEGMENTS, EACH SEGMENT DESCRIBES A FILE ON THAT DECTAPE.

E. G.



bit 0 = 1 active

DOS FILE STRUCTURE



Master File Directory (MFD)

0	-1	DUMMY WORD
1	-1	OR POINTER TO <u>BAD ALLOCATION TABLE</u> <i>Address of = 777 777</i>
2	34	SYS BLK'S FIRST BLOCK NUMBER -1 IF NOT INITIAL MFD BLOCK
3	401776	ENTRY SIZE (0-2) PLUS POINTER TO <u>STORAGE ALLOCATION TABLE</u>
4	251103	(.SIXBIT "UIC")
5	54	POINTER TO USER FILE DIRECTORY
6	400010	PROTECTION CODE FOR THIS USER (0) AND FILE DESCRIPTION ENTRY SIZE
7	0	NOT USED
10	041405	.SIXBIT "DLE"
11	60	POINTER TO UFD
12	400010	PROTECTION CODE AND FILE ENTRY SIZE
13	0	NOT USED
376	-1	POINTER TO PREVIOUS BLOCK
377	-1	POINTER TO NEXT BLOCK <i>in MFD</i>

NOTES: MFD = BLOCK #1777 if RF DISK, 47040 if RPO2

PROTECTION CODE: 1 = Protected Directory, 0 = Unprotected

ILLEGAL UFD'S: @@@, ???, and those that are current to the system - PAG, BNK, SYS, IOS.

User File Directory (UFD)

030114	.SIXBT "CALH@SRC"
100000	.
232203	.
001645	FIRST BLOCK OF THIS FILE (BIT 0 = TRUNCATION)
000002	SIZE OF FILE IN BLOCKS
001764	POINTER TO RIB BLOCK
* 200232	FILE PROTECTION CODE (0-2), START LOCATION OF RIB
142501	DATE FILE CREATED (12-21-71)
170623	.SIXBT "OFSCCLKLST"
031413	.
142324	.
002045	FIRST BLOCK OF THIS FILE
000116	SIZE OF FILE IN BLOCKS
001643	POINTER TO RIB BLOCK
* 200000	FILE PROTECTION CODE (0-2) AND START LOCATION OF RIB
142501	DATE FILE CREATED (12-21-71)

NOTES: PROTECTION CODE: (Valid only if directory is protected)
1 = Unprotected, 2 = Write Prot., 3=R/W Prot.

*RIB: The RIB may occupy its own block or, if room, occupy an area at the end of the file it is describing.

TRUNCATION: File was not closed.

↑C

DOS-15 V3A000
SA TT -12

A TT -12/RK -14
LP

SA DK -14

\$DUMP

DUMP V3A000
>1777#

1777# MFD *RK = 1777 #*

UIC-PAG

UIC=105

0	777777	777777	000034	401776	111723	000053	400010	000000
10	200107	000066	400010	000000	021613	000064	400010	000000
20	230322	001461	000010	000000	020410	000653	000010	000000
30	200523	001451	000010	000000	221310	001733	000010	000000
40	TO	367	CONTAINS 000000					
370	000000	000000	000000	000000	000000	000000	777777	777777

DUMP V3A000
>1776#

STORAGE ALLOCATION TABLE

1776# SAT *RK 1776 #*

0	004000	004000	002432	777777	777777	777777	777777	777777
10	TO	67	CONTAINS 777777					
70	777777	775777	777777	777777	777777	777777	777777	777777
100	TO	107	CONTAINS 777777					
110	777777	777777	777777	776520	000000	000000	000000	000000
120	TO	367	CONTAINS 000000					
370	000000	000000	000000	000000	000000	000000	777777	777777

DUMP V3A000
>66#

66# PAG UFD

NOT

START

DATE

0	040424	000000	021116	000663	000013	000707	200136	073104
10	053005	032524	021116	000666	000003	000672	200146	073104
20	061703	011400	021116	000674	000022	000761	200062	073104
30	561411	022200	021116	000713	000077	001327	200254	031205
40	561417	010400	021116	001356	000011	001376	200244	101204
50	111623	011414	232203	001205	000006	001217	200172	073104
60	111623	052222	232203	001221	000010	001445	200026	101204
70	111623	242203	021116	001244	000001	001244	200170	073104

100	TO	367	CONTAINS 000000					
370	000000	000000	000000	000000	000000	000000	777777	777777

DUMP V3A000
>653#

653# BDH UFD

0	062416	611516	021116	001455	000001	001455	200152	030305
10	042515	000000	021116	001457	000001	001457	200066	030305
20	042515	153100	300324	001465	000001	001465	200026	030305
30	042515	153100	300325	001467	000022	001555	200002	030305
40	561411	022265	021116	001534	000001	001534	200052	030305
50	TO	367	CONTAINS 000000					
370	000000	000000	000000	000000	000000	000000	777777	777777

DUMP V3A000
>

PATCH:

SUBJECT: DOS-15 PATCH TO DUMP V9A

The following patch corrects a problem in DUMP which outputs incorrect information on selective dumps.

<u>LOCATION</u>	<u>OLD CONTENTS</u>	<u>NEW CONTENTS</u>	<u>NEW SYMBOLIC</u>	<u>COMMENTS</u>
16256	217406	617472	JMP PATCH	/Patch area
17472	-	116034	JMS DEVICE	/Device check
17473	-	217406	LAC (-1)	/Restore inst.
17474	-	616257	JMP BACK	/Return
17224	106400	206400		

DOS-15 VIA
\$MICLOG SYS

\$DUMP

DUMP V9A
↑C

\$R PATCH

.DAT	DEVICE	UIC	USE
-14	DKA	SYS	I/O - SYS DEV
-10	TTA	SYS	SECONDARY INPUT

\$PATCH

PATCH V10A

```
>DUMP
>L 16256
>16256/217406>617472
  16257/057172>
>L 17472
>17472/215116>116034
  17473/253512>217406
  17474/055116>616257
  17475/214753>
>L 17224
>17224/106400>206400
>EXIT
```

DOS-15 VIA
\$DUMP

DUMP V9B
>

)

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)

.TITLE PRB.

FIRST PRINTING, FEBRUARY 1974

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.EJECT

/COPYRIGHT 1970, DIGITAL EQUIPMENT CORP., MAYNARD, MASS.
 /PRB.
 /PRB.#IOPS PAPER TAPE READER HANDLER == IOPS ASCII .Y.
 /M. SIFNAS/J. MURPHY
 /7-30-68

/CALLING SEQUENCE:
 /.INIT
 /CAL+.DAT SLOT(9-17)
 /1
 /0
 /0
 /.READ
 /CAL+D.M.(8-8)+.DAT SLOT (9-17)
 /10
 /LINE BUF. ADDR.
 /-WC OF L.B. (2'S COMP)
 /.WAIT
 /CAL+.DAT SLOT (9-17)
 /12

000003 A .MED=3
 700101 A RSP=700101
 700102 A RCF=700102
 700112 A RRB=700112
 700104 A RSA=700104
 700144 A RSB=700144

00000 R 040340 R	PRB.	.GLOBL PRB.	
00001 R 040341 R		DAC PRCALP	/CAL POINTER
00002 R 440341 R		DAC PRARGP	/ARG POINTER
00003 R 220341 R		ISZ PRARGP	/INDEX TO FUNCTION ADDR.
00004 R 440341 R		LAC+ PRARGP	/FUNCTION
00005 R 340422 R		ISZ PRARGP	/INDEX TO NEXT ARGUMENT.
00006 R 040014 R		TAD (JMP PRTABL	
00007 R 700314 A		DAC PRTABL	
00010 R 750100 A		IORS	/SET ION=IOF SWITCH
00011 R 777740 A		SMAICLA /AS FUNCTION OF ITS STATE	
00012 R 340160 R		LAW 17740	/ON ENTRY INTO CAL LEVEL
00013 R 040052 R		TAD PRION	
00014 R 740040 A		DAC PRDBK	
00015 R 600031 R	PRTABL	XX	
00016 R 740000 A		JMP PRIN	/1#.INIT
00017 R 600056 R		NOP	/2#.DELET. .RENAM. .FSTAT = IGNI
00020 R 600027 R		JMP PRSEEK	/3#.SEEK = IGNORED.
00021 R 600027 R		JMP PRERG	/4#.ENTER
00022 R 600045 R		JMP PRERG	/5#.CLEAR
00023 R 600052 R		JMP PRWAIT	/6#.CLOSE
00024 R 600075 R		JMP PROBK	/7#.MTAPE = IGNORED.
00025 R 600027 R		JMP PRRED	/10#.READ
00026 R 600000 R		JMP PRERG	/11#.WRITE
00027 R 760006 A	PRERG	JMP PRWATR	/12#.WAIT OR .WAITR
00030 R 620423 R		LAW 6	/ILL. FUNCTION-CAL ADDR IN .MED
		JMP+ (.MED+1	
		/INIT PTR ROUTINE	

```

00031 R 440341 R PRIN ISZ PRARGP /INDEX TO BUFFER SIZE RETURN REG
00032 R 200424 R LAC (64
00033 R 000341 R DAC+ PRARGP /STANDARD BUFFER SIZE=52
00034 R 440341 R ISZ PRARGP /INDEX TO RETURN ADDRESS
00035 R 000050 A CAL 50 /API ADDR. = ONCE ONLY CODE; THE
00036 R 000010 A PRLBHP 10 /.SETUP = L.B.H. POINTER,
00037 R 700101 A PROBP RSF / = DATA WORD POINTER
00040 R 000140 R PARER PTRINT /PARITY ERROR SWITCH
00041 R 200043 R PRUND LAC .+2 / = I/O UNDERWAY
00042 R 040035 R PTRWC DAC .-5 / = 2'S COMP WORD COUNT
00043 R 600044 R PRCHAR JMP PRSTOP / = CHAR PROCESSED
/STOP PTR ROUTINE, CLEARING I/O SWITCH
00044 R 140041 R PRSTOP DZM PRUND /CLEAR I/O UNDERWAY INDICATOR
/PTR WAIT
00045 R 200041 R PRWAIT LAC PRUND /STILL READING
00046 R 741200 A SNA
00047 R 600052 R JMP PRDBK
00050 R 200340 R PRBUSY LAC PRCALP /BUSY! RETURN
00051 R 040341 R DAC PRARGP /TO USER CAL.
00052 R 740040 A PRDBK XX /ION OF IOP
00053 R 703344 A DBR /DEBREAK FROM CAL LEVEL AND REST
00054 R 400055 R XCT .+1
00055 R 620341 R JMP+ PRARGP /EXIT
00056 R 000052 R PRNOR+PRDBK
00057 R 440341 R PRSEEK ISZ PRARGP
00060 R 600052 R PRDBKI JMP PRDBK
00061 R 761000 A PRWATR LAW 1000
00062 R 520340 R AND+ PRCALP
00063 R 741200 A SNA
00064 R 600045 R JMP PRWAIT / WAIT
00065 R 200340 R LAC PRCALP / WAITR
00066 R 500425 R AND (700000 /L,XM,MP
00067 R 040340 R DAC PRCALP
00070 R 220341 R LAC+ PRARGP /BUSY ADDRESS
00071 R 500426 R AND (77777
00072 R 240340 R XOR PRCALP
00073 R 040340 R DAC PRCALP
00074 R 440341 R ISZ PRARGP /TO NON-BUSY
00075 R 600045 R JMP PRWAIT
/PTR READ ROUTINE
00076 R 200041 R PRRED LAC PRUND /I/O UNDERWAY?
00077 R 740201 A SZAIOMA /NO START IT UP
00078 R 600050 R JMP PRBUSY /YES-WAIT IN BUSY LOOP BACK TO C
/START UP PTR
00100 R 040041 R PRSTRY DAC PRUND /SET I/O UNDERWAY SWITCH (777777)
00101 R 200057 R LAC PRDBKI /EXIT TO USER-MAINSTREAM
00102 R 040137 R DAC PTROUT
00103 R 220341 R PRNEXR LAC+ PRARGP /L.B.H POINTER IN CALL
00104 R 040037 R DAC PROBP
00105 R 040036 R DAC PRLBHP /L.B.H. POINTER
00106 R 440341 R ISZ PRARGP
00107 R 220341 R LAC+ PRARGP

```

00110	R	040042	R	DAC PTRNC	/-L.B.W.C. (2'S COMP)
00111	R	440341	R	ISZ PRARGP	/INDEX TO POINT TO EXIT
00112	R	767000	A	LAW 7000	
00113	R	520340	R	AND* PRCALP	/CHECK FOR IOPS ASCII MO
00114	R	540427	R	SAD (2000	
00115	R	000120	R	JMP .+3	
00116	R	700007	A	LAW 7	/ILLEGAL DATA MODE
00117	R	020423	R	JMP* (.MED+1	
00120	R	200430	R	LAC (1002	/IOPS ASCII DATA MODE AN
00121	R	040347	R	DAC PROTCT	/WD. PAIR CT. OF
					/1 FOR HEADER.
00122	R	100317	R	JMS PRNXWD	/INDEX PAST L.B. HEADER
00123	R	100317	R	JMS PRNXWD	/FOR IOPS ASCII
00124	R	777773	A	LAW 17773	/5/7 CHAR
00125	R	040344	R	DAC PTR57	/COUNTER
00126	R	140346	R	OZM PRCCT	/CLEAR CHAR CT.
00127	R	140345	R	OZM PRBCT	/CLEAR ASCII 8TH BIT SET COUNTER
00130	R	140040	R	OZM PARER	/CLEAR PARITY ERR. SWITCH
00131	R	700314	A	IORS	
00132	R	500431	R	AND (1000	
00133	R	740200	A	SZA	
00134	R	000172	R	JMP PREOM	
00135	R	700002	A	IOP	
00136	R	700104	A	RSA	
00137	R	740040	A	PTROUT XX	/JMP PRDBK OR JMP PRDISM
				/PTR INTERRUPT SERVICE	
00140	R	000151	R	PTRINT JMP PTRPIC	/PIC ENTRY
00141	R	040342	R	DAC PTRAC	/API ENTRY, SAVE AC
00142	R	200140	R	LAC PTRINT	/PIC OR API, L, EM, MP
00143	R	040343	R	DAC PROUT	
00144	R	700314	A	IORS	
00145	R	750100	A	SMA!CLA	
00146	R	777740	A	LAW 17740	/PIC OFF
00147	R	340160	R	TAD PRION	/PIC ON
00150	R	000155	R	JMP PRSION	
00151	R	040342	R	PTRPIC DAC PTRAC	/SAVE AC
00152	R	220432	R	LAC* (0	/PIC-PC, L, EM, MP
00153	R	040343	R	DAC PROUT	/SAVE FOR EXIT
00154	R	200160	R	LAC PRION	
00155	R	040352	R	PRSION DAC PRSW	
00156	R	700112	A	RRB	/READ PTR BUFFER
00157	R	040043	R	DAC PRCHAR	
00158	R	700042	A	PRION ION	
00161	R	200165	R	LAC PRDSMI	
00162	R	040137	R	DAC PTROUT	
00163	R	200041	R	LAC PRUND	/CHECK FOR STOP SINCE LA
00164	R	741200	A	SNA	/O.K.
00165	R	000200	R	PRDSMI JMP PRDISM	/IGNORE LAST READ, STOP
00166	R	700314	A	IORS	
00167	R	500431	R	AND (1000	/PTR NOT READY? (IORS=1)
00170	R	741200	A	SNA	
00171	R	000207	R	JMP PRIOA	

```

/END OF PAPER TAPE ROUTINE
00172 R 200433 R PREOM LAC (15 /CLEAR I/O UNDERWAY
00173 R 040043 R DAC PRCHAR /FAKE OUT END OF LINE TEST
00174 R 200423 R LAC (4 /CHANGE MODE
00175 R 340347 R TAD PRDTCT /TO EOM
00176 R 040347 R DAC PRDTCT
00177 R 000257 R JMP PRPAD
00200 R 200342 R PRDISM LAC PTRAC
00201 R 400352 R XCT PRSW
00202 R 703344 A DBR /DEBREAK FROM HANDLER LEVEL.
00203 R 400204 R XCT .+1
00204 R 020343 R JMP PROUT

/END LINE
00205 R 140041 R PRIOBB DZM PRUND /CLEAR INPUT UNDERWAY INDICATOR
00206 R 400137 R XCT PTROUT /EXIT

/PROCESS IOPS ASCII
00207 R 200042 R PRIOA LAC PTRWC /SEE IF EXCESS DATA
00210 R 740100 A SMA
00211 R 000277 R JMP PRASE3 /YES==CONTINUE UNTIL C.

/COMPUTE PARITY AND EXIT IF NULL
00212 R 777770 A LAW 17770 /PARITY COUNTER (-8)
00213 R 040350 R DAC PRCNT
00214 R 140351 R DZM PRCNT1
00215 R 200043 R LAC PRCHAR
00216 R 500434 R AND (177
00217 R 540435 R SAD (12
00220 R 000131 R JMP PROUT2 /IGNORE LF
00221 R 540436 R SAD (13
00222 R 000131 R JMP PROUT2 /IGNORE VT
00223 R 540437 R SAD (14
00224 R 000131 R JMP PROUT2 /IGNORE FF
00225 R 200043 R LAC PRCHAR
00226 R 741200 A SNA
00227 R 000131 R JMP PROUT2 /NULL
00230 R 740020 A RAR
00231 R 741400 A SZL
00232 R 440351 R ISZ PRCNT1 /1 BIT COUNTER
00233 R 440350 R ISZ PRCNT
00234 R 000230 R JMP .-4
00235 R 741400 A SZL
00236 R 440345 R ISZ PRCT /8TH BIT=1, ADD TO COUNT
00237 R 440346 R ISZ PRCT /PARITY COUNT-SHOULD BE EVEN
00240 R 200351 R PRIOA1 LAC PRCNT1
00241 R 740020 A RAR
00242 R 741400 A SZL
00243 R 440040 R ISZ PARER /NOT EVEN PARITY
00244 R 100324 R JMS PRENDT /CONVERT ALTMODES
00245 R 000434 R AND (177 /DROP ALL BUT 7 BITS
00246 R 540434 R SAD (177 /DELETE CODE (RUBOUT)-IGNORE
00247 R 000131 R JMP PROUT2
00250 R 100353 R JMS PRPK57 /PACK INTO L.B. IN 5/7
00251 R 100324 R JMS PRENDT

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00252 R 740100 A SMA
00253 R 600131 R JMP PROUT2 /NEXT ASCII CHAR
00254 R 200340 R LAC PRCCT
00255 R 540440 R SAD C1
00256 R 000124 R JMP PR1CR /IGNORE SINGLE CR LINE
00257 R 200344 R PRPAD LAC PTR57 /WORD COUNT ALL SET.
00260 R 540415 R SAD PRSCNT
00261 R 000265 R JMP PRASE
00262 R 750000 A CLA /PAD LAST
00263 R 100353 R JMS PRPK57 /WORD PAIR
00264 R 000257 R JMP PRPAD

/END OF IOPS ASCII LINE
00265 R 200347 R PRASE LAC PRDCT /WD, PAIR COUNT (INCL. HDR.)
00266 R 060030 R DAC* PRLBHP /WD.0 L.B.H.
00267 R 200345 R LAC PRDCT /DID ALL CHAR'S HAVE BIT 8
00270 R 540340 R SAD PRCCT /NO - IOPS ASCII CHECK PARITY
00271 R 600277 R JMP PRASE3 /YES - ASSUME NON IOPS ASCII
00272 R 200040 R PRASE2 LAC PARER /PARITY ERROR
00273 R 740200 A SZA /NO
00274 R 200441 R LAC C20 /YES
00275 R 200036 R PRASE4 XOR* PRLBHP /PARITY
00276 R 060030 R DAC* PRLBHP /ERROR INDICATOR.
00277 R 100324 R PRASE3 JMS PRENDT /SKIP TO END LINE
00300 R 751101 A SPAICLAICMA
00301 R 000205 R JMP PR100B /C.R. FOUND - EXIT.
00302 R 340037 R TAD PRDBP
00303 R 040037 R DAC PRDBP /POINTS TO LAST CHAR
00304 R 777400 A LAW 17400
00305 R 520037 R AND* PRDBP
00306 R 240442 R XOR C33 /PUT CR IN LAST WORD PAIR
00307 R 060037 R DAC* PRDBP
00310 R 440037 R ISZ PRDBP /INCASE MORE BEFORE CR
00311 R 220036 R LAC* PRLBHP
00312 R 500443 R AND C00
00313 R 740200 A SZA
00314 R 000131 R JMP PROUT2 /VALIDITY BITS ALREADY S
00315 R 200443 R LAC C00 /LINE BUFFER OVERFLOW.
00316 R 000275 R JMP PRASE4
00317 R 000000 A PRNXWD 0
00320 R 440037 R ISZ PRDBP /INDEX TO NEXT DATA WORD
00321 R 440042 R ISZ PTRWC /INDEX WORD COUNT
00322 R 020317 R JMP* PRNXWD /EXIT FOR NEXT CHAR
00323 R 000265 R JMP PRASE /EXIT TO END OF IOPS ASCII LINE

/END LINE TEST - CONVERTS ALTMODE TO STANDARD 175
PRENDT 0
00324 R 000000 A LAC PRCHAR
00325 R 200043 R AND C177
00326 R 500434 R SAD C15 /RETURN
00327 R 540433 R LAW 15
00330 R 700015 A SAD C175 /ALTMODE
00331 R 540444 R LAW 175
00332 R 700175 A SAD C176 /ALTMODE
00333 R 540445 R

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00334 R 760175 A LAW 175
00335 R 540442 R SAD (33 /ESCAPE
00336 R 760175 A LAW 175
00337 R 620324 R JMP* PRENDT
/VARIABLES=NOT SAVED=APPLY TO CURRENT ACTIVE REQUEST
00340 R 000000 A PRCALP 0 /CAL POINTER
00341 R 000000 A PRARGP 0 /ARG. LIST AND EXIT POINTER
00342 R 000000 A PTRAC 0 /SAVED AC(INTERRUPT)
00343 R 000000 A PROUT 0 /PC,L,EM,MP
00344 R 000000 A PTR57 0 /CHAR. POSITION COUNTER IN 5/7 PAIR
00345 R 000000 A PROCT 0 /ASCII=WITH=8TH=BIT=SET=CHAR COUNTER
00346 R 000000 A PRCCT 0 /CHAR CT.
00347 R 000000 A PROTCT 0 /DATA WORD PAIR IN LINE COUNTER
00350 R 000000 A PRCNT 0 /PARITY CHECK COUNTER
00351 R 000000 A PRCNT1 0 /1 BIT COUNTER FOR PARITY CHECK
00352 R 000000 A PR5W 0 /ION OR IOP
/5/7 IOPS ASCII PACKING ROUTINE.
/ PTR57 IS INITIALIZED TO 777773
/ PRIOR TO THE 1ST CALL.
/
00353 R 000000 A PRPK57 0 /CHAR. IN AC BITS 11-17.
00354 R 742020 A RTR /MOVE TO AC BITS 0-6
00355 R 742020 A RTR
00356 R 742020 A RTR
00357 R 742020 A RTR
00360 R 040350 R DAC PRTMP
00361 R 777771 A LAW 17771 /-7
00362 R 040351 R DAC PRLPCT
00363 R 200350 R PRPKBK LAC PRTMP /ROTATE CHAR LEFT
00364 R 740010 A RAL /7 BITS THROUGH
00365 R 040350 R DAC PRTMP /THE DOUBLE WORD
00366 R 200421 R PRBCK2 LAC PRRTHF /ACCUMULATOR
00367 R 740010 A RAL /PRLPHF/PRRTHF.
00370 R 040421 R DAC PRRTHF
00371 R 200420 R LAC PRLFHF
00372 R 740010 A RAL
00373 R 040420 R DAC PRLFHF
00374 R 200351 R LAC PRLPCT
00375 R 745200 A SNA!CLL
00376 R 000404 R JMP PRPDNE /2 WORDS ALL SET.
00377 R 440351 R ISZ PRLPCT /IS 7 TIMES COUNT EXHAUSTED?
00400 R 000363 R JMP PRPKBK /NO.
00401 R 440344 R ISZ PTR57 /DO WE HAVE 5 CHARS.
00402 R 620353 R JMP* PRPK57 /NO. EXIT
00403 R 000366 R JMP PRBCK2 /SHIFT LEFT ONCE MORE!
00404 R 200420 R PRPDNE LAC PRLFHF /PLACE ACCUMULATED
00405 R 060037 R DAC* PRDBP /2 WORDS INTO
00406 R 100317 R JMS PRNXWD /USERS LINE BUFFER,
00407 R 200421 R LAC PRRTHF /UPDATING POINTERS.
00410 R 060037 R DAC* PRDBP
00411 R 200347 R LAC PROTCT /INCREMENT
00412 R 340431 R TAD (1000 /DATA WD. PAIR

```

```

00413 R 040347 R DAC PRDTCT /COUNT
00414 R 100317 R JMS PRXWD
00415 R 777773 A PRSCNT /RESET 5 CHAR COUNTER
00416 R 040344 R DAC PTR57
00417 R 020353 R JMP* PRPK57
      000350 R PRTMP=PRCNT /TEMP. STORAGE FOR 5/7 CHAR.
      000351 R PRLPCT=PRCNT1 /ROTATE 7 BITS COUNTER.
00420 R 000000 A PRLFHF 0 /2 WORD ACCUMULATOR FOR
00421 R 000000 A PRRTHF 0 /5/7 WORD PAIR,
      000000 A .END
00422 R 000014 R *L
00423 R 000004 A *L
00424 R 000004 A *L
00425 R 700000 A *L
00426 R 077777 A *L
00427 R 002000 A *L
00430 R 001002 A *L
00431 R 001000 A *L
00432 R 000000 A *L
00433 R 000015 A *L
00434 R 000177 A *L
00435 R 000012 A *L
00436 R 000013 A *L
00437 R 000014 A *L
00440 R 000001 A *L
00441 R 000020 A *L
00442 R 000033 A *L
00443 R 000000 A *L
00444 R 000175 A *L
00445 R 000176 A *L

```

SIZE=00446 NO ERROR LINES

TABLE OF CONTENTS

<u>SECTION</u>	<u>PAGE NO.</u>
INTRODUCTION	154
UNICHANNEL 15 HARDWARE ARCHITECTURE	155
UNICHANNEL 15 SOFTWARE ARCHITECTURE	158
MEMORY LAYOUT	162
PIREX TASKS	163
INTERRUPT LINK	167
MX15-B MEMORY MULTIPLEXERS	172
PDP-11/05 CONTROL REGISTERS	171
SYSTEM CONFIGURATION	174
SYSTEM RESTRICTIONS	175
PDP-15 UNICHANNEL OPTIONS	177

INTRODUCTION

This guide describes in more detail, the UNICHANNEL 15 operation and features presented in the RK15 Disk Cartridge System Option Bulletin.

The first sectionpresents a look at the UC15 system architecture.

The second section.....describes the PIREX monitor system; how to use it and other software aids.

The final section.....provides, for those interested in creating their own programs, a complete hardware specification including IOT and register descriptions.

Supplementing this guide are two manuals:

Unichannel 15 System Maintenance Manual:..DEC-15-HUCMA-B-D

UC15 Software Manual:.....DEC-15-XUCMA-A-D

The maintenance manual describes the details of the MX15-B and the DR15-C logic and gives maintenance details.

The software manual describes the details of the PIREX Monitor.

UNICHANNEL - 15 HARDWARE ARCHITECTURE

The term UNICHANNEL was created because it emphasizes the union of Digital's UNIBUS with the big computer concept of the programmable I/O channel. UNICHANNEL 15 unites low cost, mass produced peripherals with big computer software and performance on the PDP-15.

UNICHANNEL 15 (UC15) is a peripheral processor for the PDP-15 utilizing the PDP-11/05 minicomputer. It provides the PDP-15 with a second general purpose processor and a second high speed I/O bus; the UNIBUS. This UNIBUS is an 18-bit pathway permitting transfer of either 18-bit words, 16-bit PDP-11 words, or two 8-bit bytes.

The UC15 allows flexible low cost configuration and expansion of PDP-15 systems.

The UC15 minimizes the peripheral processing load on the PDP-15 allowing maximum computational throughput in a low-priced, medium scale system.

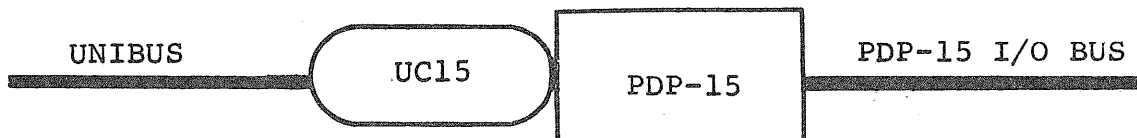


FIGURE 1: Simplified UC15 Diagram

UNICHANNEL 15 OPERATION

There are three major components of the UC15:

1. A PDP-11/05 computer with "local" PDP-11 memory.
2. An MX15-B memory multiplexer which allows both the PDP-15 processor and the PDP-11 processor to share common memory. The shared memory is ordinary 18-bit PDP-15 core memory.
3. An "interrupt link" to provide a real-time means of inter-processor communications.

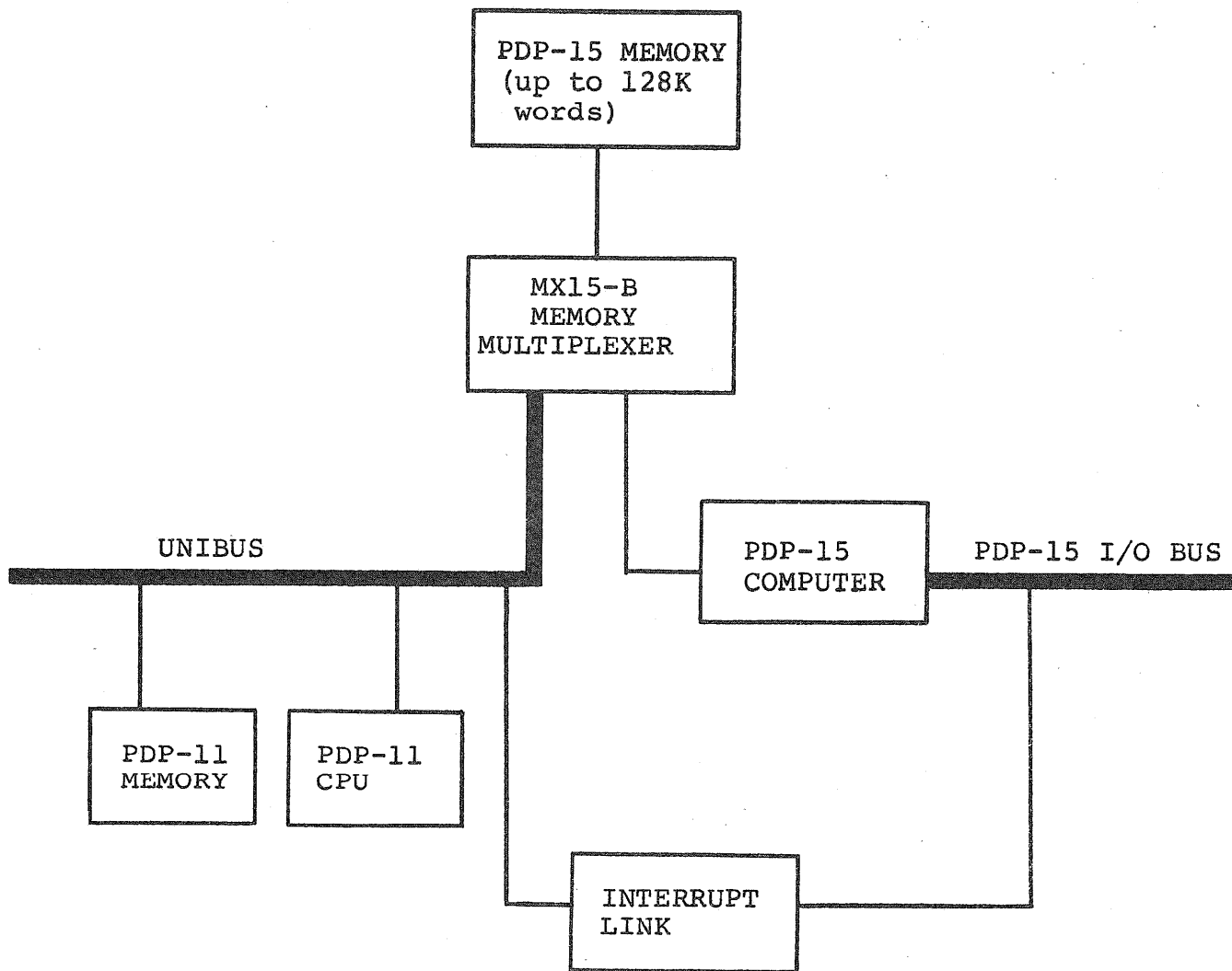


Figure 2: Diagram of UC15 Hardware Interrelationships

SUMMARY - UNICHANNEL 15 HARDWARE ARCHITECTURE

This particular architecture was chosen because of its many advantages.....

PDP-15 Memory is addressable by the UNIBUS. Hence, DMA transfers from and to such secondary storage devices as disks are direct.

The interrupt link provides inter-processor signaling on a micro-second basis. This is ideal for efficient real-time service -- a necessity for flexible I/O control.

All PDP-15 systems may be upgraded by adding the UC15. All memory remains useable.

Cost is minimized by allowing the PDP-11 to share the PDP-15 console and paper tape loader system.

Maximum use of the PDP-15 memory is maintained through synchronization overlap with memory use by the MX15-B. This "pre set up" technique increases the number of memory cycles per second when both PDP-15 and PDP-11/05 are accessing the common PDP-15 memory.

The UNIBUS provided by the UC15 is electrically compatible with any device meeting UNIBUS interfacing specifications with the following restraints:

1. UNIBUS lengths must be kept short.
2. No provision is made for UNIBUS parity.

Data in the common PDP-15 memory may be treated as either 18 or 16 bit words or as (2) 8-bit bytes.

True simultaneous parallel processing is possible in the local and common memories.

The DMA rate is high and the worst case and average latencies are low for maximum I/O performance.

Finally, the system is highly modular allowing flexibility in configuration and excellent software utilization and control. The system permits variations in both local and common memory size. It allows almost any combination of PDP-15 and UNIBUS peripherals.

UNICHANNEL - 15 SOFTWARE ARCHITECTURE

The hardware architecture is complimented by sophisticated system software. PDP-15 software systems running with a UNICHANNEL system relies on PIREX, a compact multitasking peripheral executive. In addition to PIREX, Digital supplies UNIBUS device drivers, UNICHANNEL compalible handlers, and supporting utility functions.

The software system used by UC15 consists of two parts:

1. One component is a mutli-programming peripheral processor executive called PIREX and is executed by the PDP-11.
2. The other component is an operating system in a PDP-15. (e.g. DOS-15 or BOSS-15).

PIREX

PIREX is a multi-programming executive designed to accept any number of requests from a PDP-15 or PDP-11 and process them on a priority basis while processing other tasks concurrently. PIREX services all Input/Output requests from the 15 in parallel on a controlled priority basis. Requests to busy routines (called tasks) are automatically queued (entered into a waiting list) and processed whenever the task in reference is free. In a background environment, PIREX is also capable of supporting any number of priority driven software tasks initiated by the 15 or the 11 itself.

Figure 3 shows the communications flow in a UNICHANNEL system. The possible links which may exist in the system are as follows:

1. Handler to driver to allow the PDP-15 to use a UNICHANNEL device.
2. Handler to non-driver task to allow the PDP-11 to intercept output and manipulate it or store it for spooling.
3. Program to non-driver task to allow cooperative processing on the two CPU's as occurs in the use of the MAC-11 assembler.

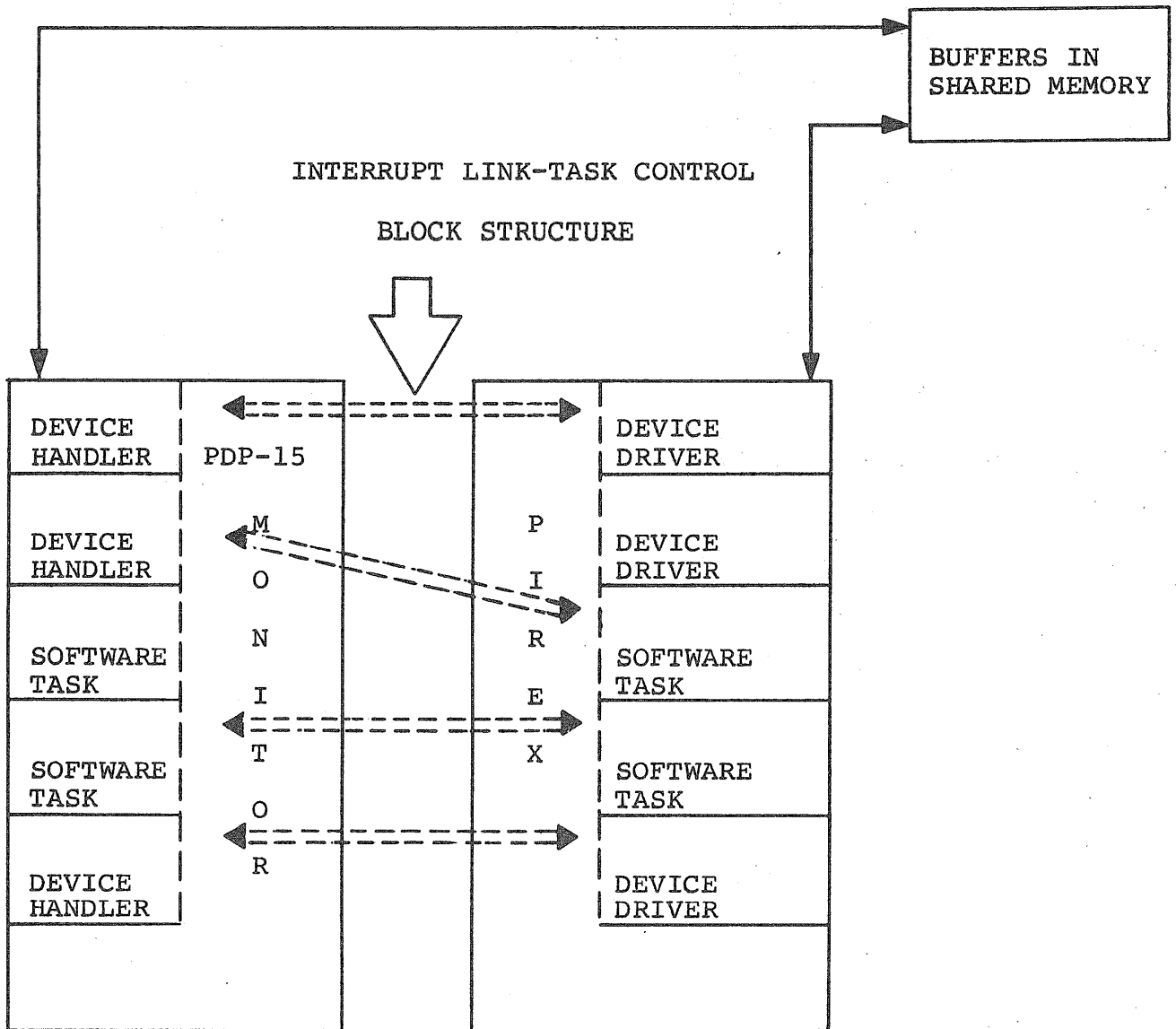


Figure 3 : DIAGRAM OF UNICHANNEL SOFTWARE SYSTEM

UNICHANNEL ADDRESSING

The Unichannel system makes use of a PDP-11 as an intelligent peripheral controller for the larger PDP-15 main computer. In order to effectively operate with a minimum of interference with the PDP-15, the PDP-11 uses its own LOCAL MEMORY of between 4K and 12K 16-bit words.

COMMON MEMORY is that memory directly accessible to both the PDP-15 and the PDP-11.

COMMON MEMORY occupies the upper portion of the PDP-11 address space and, at the same time, the lower portion of the PDP-15 address space.

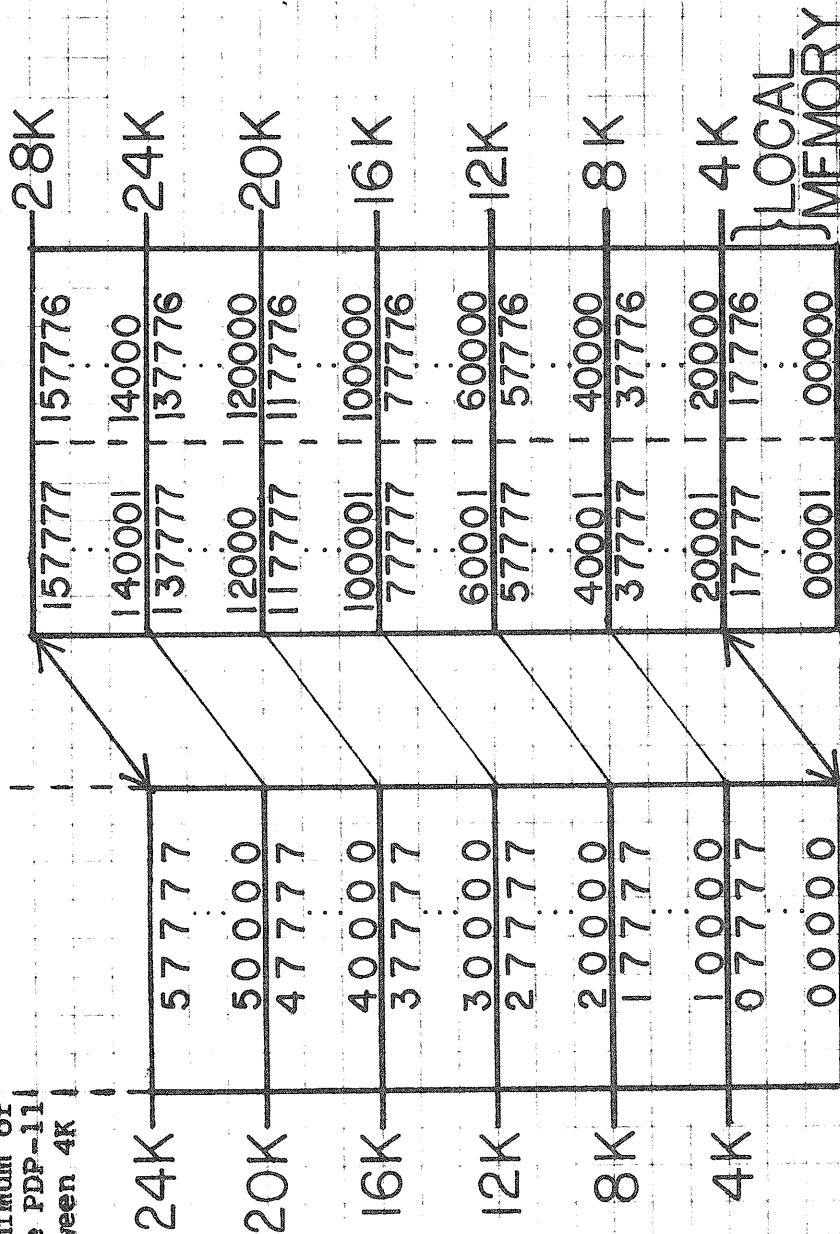
NOTE: PDP-11 LOCAL MEMORY

PDP-15/PDP-11 COMMON MEMORY

must not exceed 28K. DOS-15 requires a minimum of 16K of memory on the PDP-15. Therefore, the PDP-11 LOCAL MEMORY may be 12K or 8K or 4K.

NOTE: The PDP-11 is a byte oriented machine. The 8-bit bytes are numbered sequentially with two 8-bit bytes corresponding to one 16-bit word, thus--

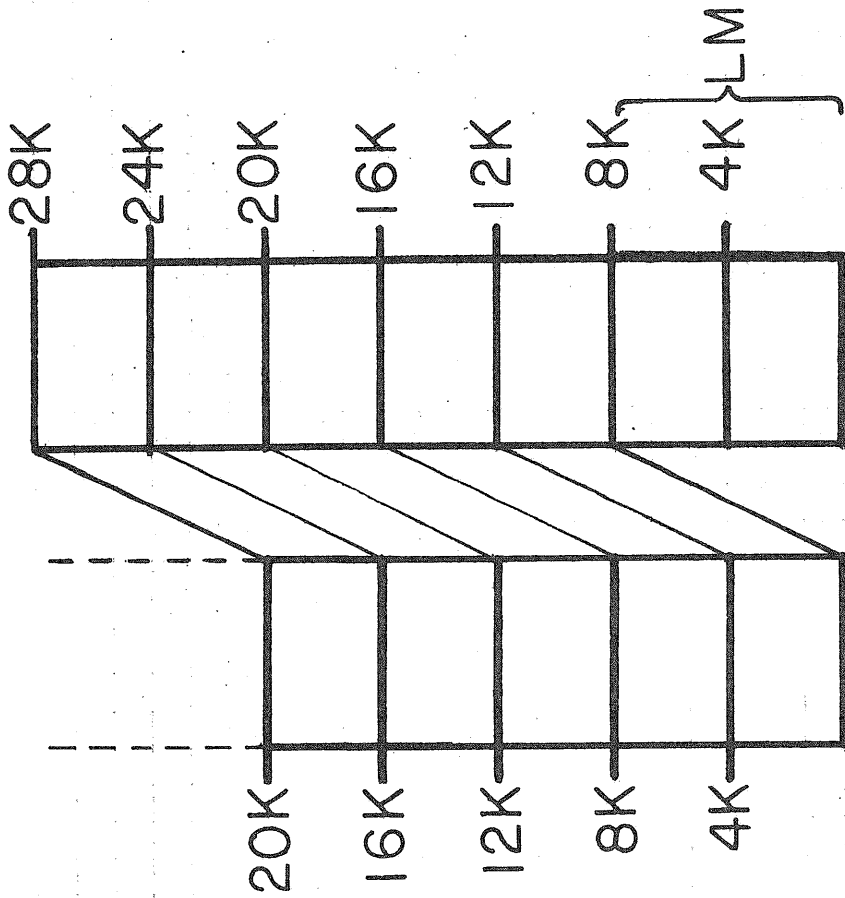
WORDS	BYTES	OCTAL ADDRESSES
4K =	8K	00000-17777
8K =	16K	00000-37777
12K =	24K	00000-57777



PDP-15

PDP-11

4K LOCAL MEMORY
24K COMMON MEMORY



PDP-15 PDP-11

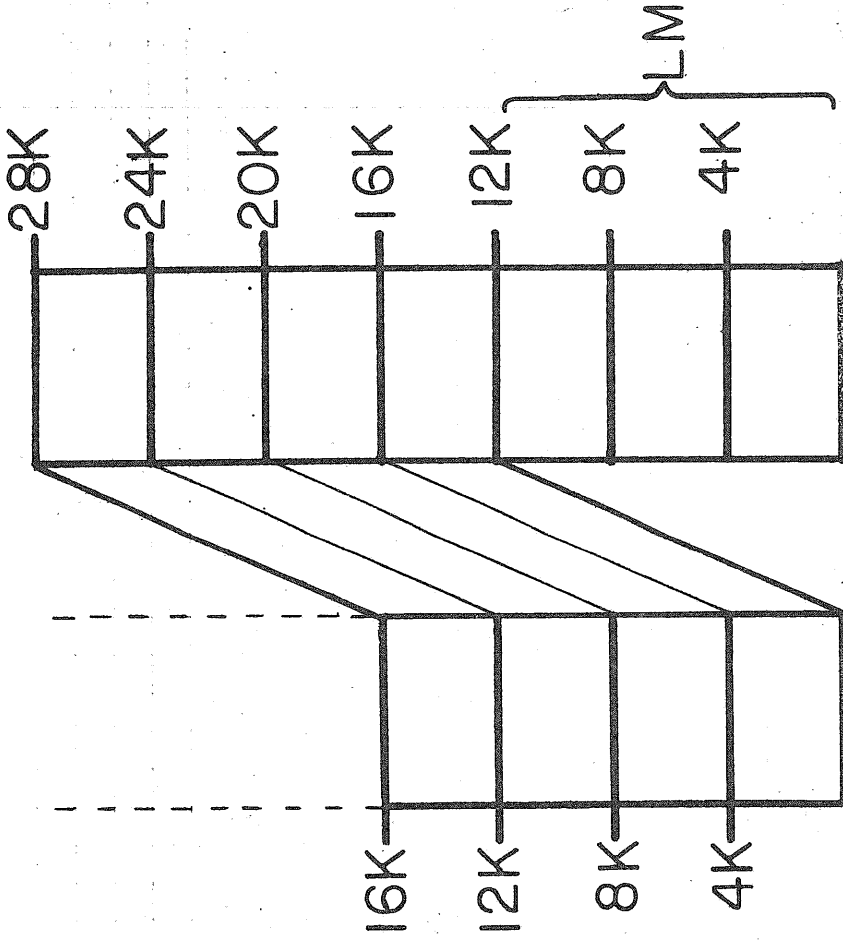
8K LOCAL MEMORY
20K COMMON MEMORY

ADDRESS CORRESPONDENCE

$AD11 = (AD15 * 2) + \text{LOCAL MEMORY size in bytes}$

$AD15 = (AD11 - \text{LOCAL MEMORY size in bytes}) / 2$

e.g. Address 7777 on the PDP-15 is address 37776 on the PDP-11 with 4K local memory
Address 60000 on the PDP-11 with 4K local memory is address 20000 on the PDP-15



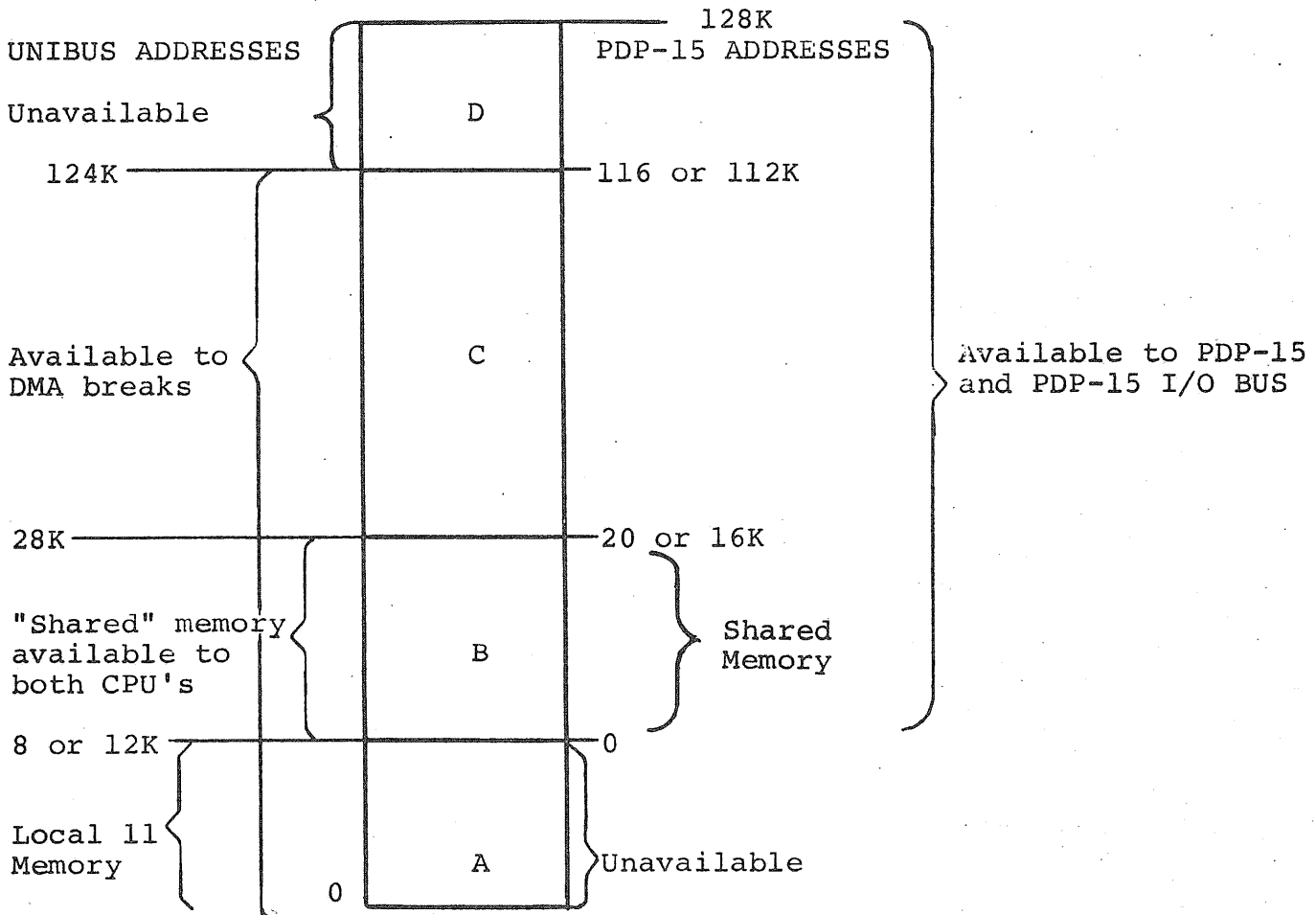
PDP-15 PDP-11

12K LOCAL MEMORY
16K COMMON MEMORY

LM: LOCAL MEMORY

MEMORY LAYOUT

Figure 4 details the memory map which exists on UNICHANNEL System. Note that both the 11 and 15 parts of the system can operate concurrently, all memory contention is resolved by the MX15-B. Note also, that if the 11 system operates with area "A" complete simultaneity is possible because no memory contention can occur.



"Local" PDP-11 Memory = A
 "Shared" Memory = B
 PDP-11 CPU Address Space = A + B
 UNIBUS DMA Address Space = A + B + C
 PDP-15 Address Space = B + C + D

Figure 4 : UNICHANNEL SYSTEM MEMORY MAP

PIREX TASKS

The PIREX software system consists of several routines to support multi-programming among tasks. These routines perform such functions as: context switching, node manipulation and scheduling. The tasks which execute in this environment are device drivers, directives to PIREX, or merely software routines which execute in a background mode.

Device drivers are tasks which typically perform rudimentary device functions (e.g.: read, write, search, process interrupts, etc.), Directives are tasks which perform some specific operation for a task under PIREX. The connecting and disconnecting tasks to/from PIREX are performed by the CONNECT and DISCONNECT directives. The third type of tasks are software routines which execute in a background mode of operation. The MACRO-11 assembler and Spooler are both run as background tasks.

To support multiprogramming among tasks, each task is required to have a format as shown in the figure below:

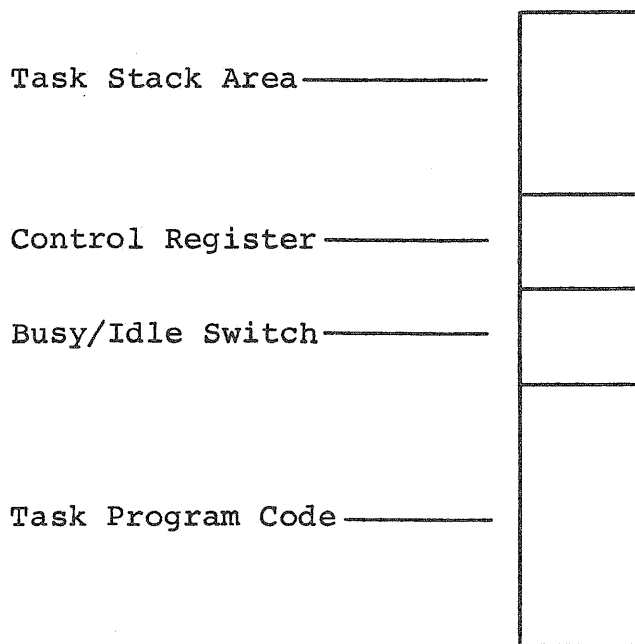


Figure 5 TASK FORMAT

The execution of a Task by PIREX is accomplished by first scanning the Active Task List (ATL). The ATL is a priority-ordered linked list of all active Tasks in the current system currently capable of running. An Active Task is one which:

1. Is currently executing.
2. Has a new request pending in its deque (double ended queue).
3. Has been interrupted by a higher priority task.

When a runnable task is found, the stack area and general purpose registers belonging to the task are restored and program control transferred to it. Program execution begins at the first location of the task program code (See Figure 2.1) or at the point where the task was previously interrupted by a higher priority task. When a task is interrupted by other tasks, its general purpose registers and stack are saved. The ATL is rescanned when a new request is issued to a task or when a previous request is complete.

When the PIREX Software System is running, it is normally executing the NUL task (a PDP-11 WAIT Instruction); The NUL task is run whenever there are no requests pending, a task suspends itself in a wait state, or while all other tasks are waiting for I/O previously initiated. When the PDP-15 issues a request to the PDP-11 to be carried out by PIREX, it does so by interrupting the 11 at Level 7 (the highest PDP-11 Interrupt Level) and simultaneously passing it an address of a Task Control Block (TCB) through the interrupt Link.

An 11 task can issue requests via the IREQ MACRO. The contents of the TCB completely describe the request (task addressed, function, optional interrupt return address and level, status words, etc....) The TCB will usually reside in the PDP-15 memory and must be directly addressable by the 11. (i.e. It resides in shared memory).

Error conditions are passed back to the 15 in the Task Control Block (TCB) along with status information necessary for complete control and monitoring of a particular request. Usually the request is to a device on the 11 but other types are allowed.

Task Control Blocks are used for communication with PIREX and tasks running under it. The general format of a TCB consists of three words followed by optional words necessary for task communication. Optional words, generally are used to pass buffer addresses, commands and device status as may be appropriate.

TCB: (API TRAP ADDRESS *400(8)) + API LEVEL
(FUNCTION CODE *400(8)) + TASK CODE NUMBER

REV: REQUEST EVENT VARIABLE
(Optional Words)

Figure 6 STANDARD TCB FORMAT

The "TRAP ADDRESS" is a PDP-15 API trap vector and has a value between 0 and 377 (8). Location 0 here corresponds to location 0 in the PDP-15. The API Level is the priority level at which the interrupt will occur in the PDP-15 and has a value between 0 and 3. A 0 signifies API "Level" 0, a 1 for level 1 etc... The API trap address and level are used by tasks in the PDP-11 when informing the 15 that the requested operation is complete (e.g...a disk block transferred or line printed).

The Task code number is a positive number between 0 and 128 that tells PIREX which task is being referenced, (Tasks are addressed by a numeric value rather than by name).

The Function Code determining whether hardware interrupts are to be used at the completion of the request. If the code has a value of 0, an interrupt is generated at completion of the request; If a 1, an interrupt is not made.

The Request Event Variable, commonly called REV or just EV, is initially cleared by PIREX (set to zero) and then set to a value "n" (by the associated task) at the completion of the request. The values of "n" are:

- 0 = request pending or not yet completed.
- 1 = request successfully completed.
- 2 = (mod 2¹⁶-1) non-existent task referenced.
- 3 = (mod 2¹⁶-1) illegal API level given (illegal values are changed to level 3 and processed).
- 4 = (mod 2¹⁶-1) illegal directive code given.
- 777 = (mod 2¹⁶-1) request node was not available from the Pool, i.e. the POOL was empty, and the referenced task was currently busy or the task did not have an ATL node in the Active Task List.

NOTE -- the Task Control Block specification clearly defines a modular communications structure with minimum impact on PDP-15 software.

ADDING DRIVERS TO PIREX

A powerful feature allows the PDP-15 to bring in a PDP-11 driver, (into either its own memory or the 11's local memory) connect it to PIREX via a connect directive (a disconnect directive) is also provided) and then issue I/O requests through PIREX to the driver. The user can now take full advantage of the existing and future PDP-11 peripherals along with an elaborate queuing structure built into PIREX allowing complete parallel processing.

MACRO 11 ASSEMBLER (MAC11) AVAILABLE)

A MACRO 11 Assembler is provided. This assembler is a Macro subset of the existing PDP-11 Macro assembler and is slightly modified to run under the control of DOS-15 and PIREX.

To accomplish this, the MACRO assembler (MAC11) is loaded by the 15 as a core image into bank 1 of the 15. MAC 11 is then connected up as a low priority driver to PIREX and requested to begin the assembly. The 11 then carries out the actual assembly while the 15 handles all of the opening and closing of files, reading and writing of test and object information until the assembly is complete. To the user at the console teletype, MAC 11 appears to be just a DOS-15 system program which is loaded in and run by the 15.

NOTE: That any customer developed software should of course, take into account PIREX context switch, the bandwidth of the UNIBUS 18 and latency consideration of the associated system.

SUMMARY

As one can easily see, the UC15 software system is a powerful tool to the user who requires the utmost in flexibility and utility. UC15 also provides an expansion capability beyond any system currently available.

INTERRUPT LINK

The following section describes the registers and control of the interrupt link. This link is used to pass Task Control Block Pointers (and through them the information in Task Control Blocks) between the PDP-15 and PDP-11 systems. The hardware which comprises this link consists of a DR15-C special purpose interface to the PDP-15, I/O BUS, and 2 DR11-C general purpose UNIBUS interfaces. The DR15-C is controlled by PDP-15 IOT's while the DR11's are accessed as registers on the UNIBUS.

Register Descriptions (PDP-11)

(CSR) 767770 Bit 6 - when bit 6 is a 1, it will enable an interrupt on BR5 to TV 300, if the API DONE flag is set in bit 7 of 767770.

Bit 7 - API DONE - set to 1 whenever none of the 4 API channels has a request pending.

NOTE: Neither of these bits is expected to be used in normal systems programming.

(ODB) 767772 Low byte - contains the API address for an API level \emptyset break. Loading a new value in the byte causes the appropriate API flag to be set in the DR15-C and an API break in the PDP-15 will occur, if the API is enabled and no higher activity is occurring. It also will cause a PI interrupt if API is not installed.

High byte - contains the API address for an API level 1 break. Same conditions as low byte.

(IDB) 767774 Bit \emptyset - contains bit "2" of the Task Control Block Pointer (TCBP). See note under bit 1.

Bit 1 - contains bit "1" of the TCBP.

NOTE: That reading 767774 does not effect the new TCBP flag in bit 7 of 767760.

Bit 6 - API 2 DONE flag - when a 1 indicates that there is no API level 2 request pending before the PDP-15. When a 1 also indicates the 767762 low byte may be loaded with a new API level 2 address to cause a new API interrupt level 2 and set the API 2 flag in the DR15-C.

Bit 7 - API \emptyset DONE flag - when a 1 indicates that there is no API level \emptyset request pending before the PDP-15. When a 1 also indicates that 767772 low byte may be loaded with a new API level \emptyset and set the API \emptyset flag in the DR15-C.

Bit 8 - Local Memory Size bit 0 - the least significant bit of a two bit field which specifies the number of 4K word memory banks that are connected to the UNIBUS.

Bit 9 - Local Memory Size Bit 1 - the most significant bit of a two bit field which specifies that number of 4K memory banks are connected to the UNIBUS.

<u>LMS1</u>	<u>LMS0</u>	
0	0	0 Local Memory
0	1	4K Local Memory
1	0	8K Local Memory
1	1	12K Local Memory

Bit 14 - API 3 DONE flag - when a 1 indicates that there is no API level 3 request pending before the PDP-15. When a 1 also indicates that 767762 high byte may be located with a new API level 1 address to cause a new API interrupt at level 3 and set the API 3 flag in the DR15-C.

Bit 15 - API 1 DONE flag - when a 1 indicates that there is no API level request pending before the PDP-15. When a 1 also indicates that 767772 high byte may be loaded with a new API level 1 address to cause a new API interrupt at level 1 and set the API in the DR15-C.

(CSR) 767760 Bit 6 - ENABLE TCBP (Task Control Block Pointer) INTERRUPT - When a 1 allows and interrupt on BR level 7 to TV 310 upon receipt of a new TCBP from the PDP-15.

Bit 7 - NEW TCBP flag - is set to 1 whenever the PDP-15 issues IOT 706006 thus placing a new TCBP in 767764 and bits 0 and 1 of 767774. It is cleared by the PDP-11 doing a DATI to location 767764.

(ODB) 767762 Low byte - contains the API address for an API level 2 break. Same conditions as 767772 low byte.

High byte - contains the API Address for an API level 3 break. Same conditions as 767772.

(IDB) 767764 TCBP (Task Control Block Pointer) - bits 3-17. This contains the lowest 15 bits of the address sent by the PDP-15. Note: that the address is "word" aligned. Note also that doing a DATI to this register lowers the New TCBP flag (767760 bit 7) and also sets the DONE flag cleared by IOT 706002 in the PDP-15.

PDP-15 IOT's

706001 SIOA - Skip I/O Accepted. Tests whether the TCBP DONE flag is set indicating the PDP-11 has read the TCBP and skips the next location if the DONE flag is a 1.

706002 CIOD - Clear I/O Done. Clear the TCBP DONE flag.

706006 LIOR - Load I/O Register and clear TCBP DONE flag. Places the contents of the PDP-15 "AC" into an 18-bit buffer register. The output of the buffer register is seen by the PDP-11 as TCBP at location 767764 and bits 0 and 1 767764. The IOT also causes the TCBP DONE flag to be cleared and in the PDP-11 causes bit 7 to be set in location 767760, which in turn causes the PDP-11 to do an interrupt at BR 7 to TV location 310.

706112 RDRS - Read Status Register - Clears the AC and loads the contents of the DR15-C status register into the AC. (This effectively moves the DR15-C enable interrupt bit into bit 17 of the AC).

706122 LDRS - Load Status Register. Loads the contents of the AC into the DR15-C status register. (Places value of AB bit 17 in the DR15-C "enable interrupts" bit).

706104 CAPI0 - Clear API0 flag in DR15-C.

706124 CAPI1 - Clear API1 flag in DR15-C.

706144 CAPI2 - Clear API2 flag in DR15-C.

706164 CAPI3 - Clear API3 flag in DR15-C.

706101 SAPIO - Tests the API0 flag in the DR15-C and skips the next instruction if the flag is 1.

- 706121 SAPI1 - Tests the API1 flag in the DR15-C and skips the next instruction if the flag is 1.
- 706141 SAPI2 - Tests the API2 flag in the DR15-C and skips the next instruction if the flag is 1.
- 706161 SAPI3 - Tests the API 13 flag in the DR15-C and skips the next instruction if the flag is 1.

PDP-15 STATUS REGISTER (DR15-C)

Bit 17 Enable PI/API interrupts. When a 1 enables interrupts from the PDP-11 processor. Note this bit is set to a 1 by initialize and the CAF instruction. It can only be cleared by using the LDRS (IOT 706122) instruction.

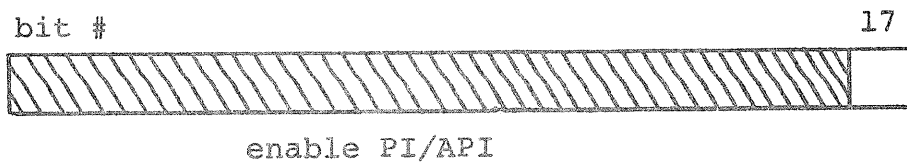
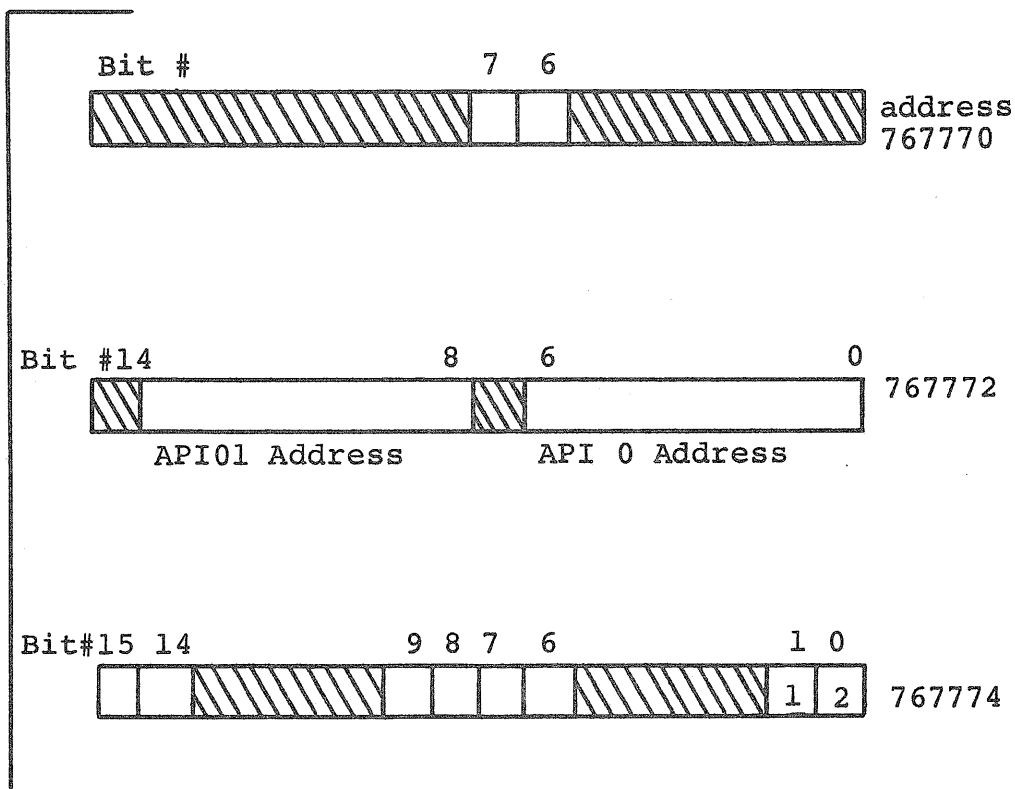


Figure 7

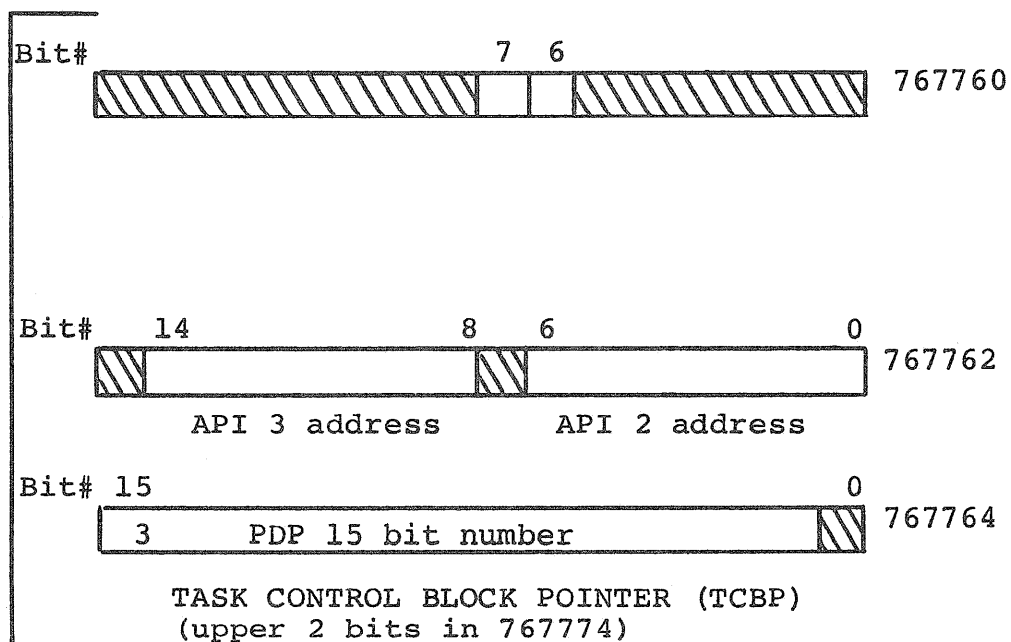
Figure: 8

PDP-11/05 CONTROL REGISTERS

DR11-C #0
TV = 300
BR = 5



DR11-C #1
TV = 310
BR = 7



MX15-B MEMORY MULTIPLEXERS

When the PDP-15 memory is accessed by the PDP-11/05 or any NPR UNIBUS device, the addresses are relocated by the MX15-B multiplexer.

The MX15-B multiplexer not only relocates the UNIBUS addresses but emulates byte operations in PDP-15 memory. Hence normal PDP-11 programs, with byte read and byte write operations may be executed from PDP-15 memory. Also such byte oriented NPR devices as Mag Tape may make transfers directly to PDP-15 memory.

Note: That the PDP-11 processor can access the PDP-15 memory which is between the end of local memory and the 28K of address space available to its address scheme.

A. Output - PDP-15 Memory Bus

Will connect to MM15, MX15-A, and ME15 memories.

B. Inputs

PDP-11: Modified UNIBUS with PA and PB used as D16 and D17 respectively. It meets all other UNIBUS specs. Defined as UNIBUS/18, input would have a lower address bound that could be fixed to any 4K multiple address 0-120K. This would be specified as jumpers. Note that only 8K and 12K of local memory will be supported by diagnostics and systems programs. Hence, the maximum commonly addressable memory (11 processor) will be 20K or 16K. An upper limit would be provided as 124K.

The addresses presented from the PDP-11 are relocated to prevent location 0 being the same physical address on each machine. The PDP-11 will be able to be relocated by 4K increments to 124K. Local PDP-11 memory is restricted to 4 increments.

Note that any "write" operation to a common memory location by 8 bit or 16 bit UNIBUS devices causes PDP-15 data bits 0 and 1 of the location to be forced 0.

PDP-15: Standard 15 Memory Bus Interface - no upper and lower bounds. No relocation. Emphasis is on minimum delay through multiplexer for this port.

SYSTEM CONFIGURATION

The UC15 cabinet will replace the current disk cabinet immediately to the left of the PDP-15 processor.

The increased spacing will require longer I/O or memory bus cables in some installations.

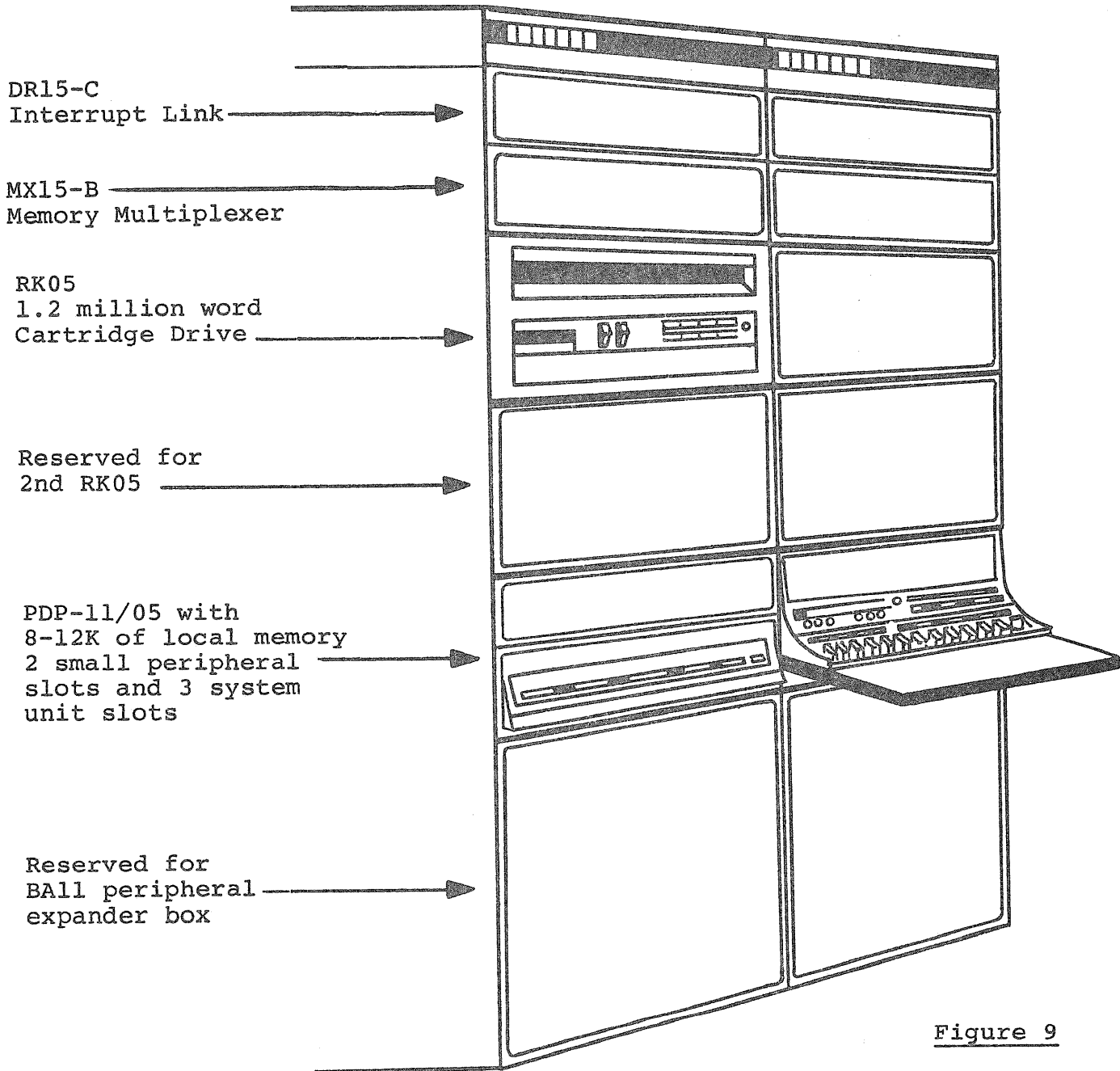


Figure 9

SYSTEM RESTRICTIONS

RK05 (RK11) Disk Pack Capability

The 18 bit RK11 disk pack will not be able to be read by RK11-C or RK11-D system (16-bit only systems).

This means that data bases and PDP-11 files created on 18-bit RK11 systems may not be taken directly to an PDP-11 only system. The transfer medium for such a transfer would have to be Mag Tape.

This situation was chosen to make RK11-C and RK11-D packs compatible (i.e.all PDP-11 only systems).

Memory Limits

UNIBUS NPR devices can access a maximum of 124K. The amount of shared memory available to UNIBUS NPR devices is 124K less the amount of local memory. In a "normal" configuration the PDP-11/05 would have 8K of memory, in which case the available PDP-15 memory would be limited to 116K. This limit is due to the fact that UNIBUS/18 peripherals must have access to all memory. The maximum memory of the 11 without some relocation option would be 28K.

Note: That the PDP-11 with 8K of local memory can only address the lowest 20K of common memory to access Task Control Blocks set up by the PDP-15.

I/O Latency

Multiport memories always have increased worst case latency over a single port-non-competitive situation. This system is no exception. The PDP-11 normally gives an "NPR break" a worst case latency to BSSY of 7.0 usec. On this system, we must add to that time, the time it requires the PDP-15 to do three I/O memory cycles (5.0 usec.). The worst case latency is, hence, 12.0 usec.

CAF/RESET Limitations

The following timing considerations are of interest to programmers:

A RESET instruction may cause the PDP-15 to incorrectly read the API address. The Console RESET and CAF instruction may violate UNIBUS specifications. Hence, random "initialize" pulses may cause system malfunctions. The following guidelines must always be followed:

1. CAF must not be executed while there is a Task Control Block Pointer (TCBP) waiting to be read by the PDP-11.
2. RESET must not be executed while there are API requests pending for the PDP-15.
3. RESET must not be executed if there is any NPR activity on the UNIBUS. All active NPR devices must be shut down in a power fail sequence prior to executing RESET.

PDP-15 UNICHANNEL OPTIONS

UC15-HE	Peripheral Processor: 11/05 or 11/10-NC or - SA, 2 DR11-C, DR15-C, MX15-B, DD11-B, KY11-JH, H950, 115V.	8K Local Memory
UC15-HF	Peripheral Processor: 11/05 or 11/10-ND or - SB, 2 DR11-C, DR15-C, MX15-B, DD11-B, KY11-JH, H950, 230V.	8K Local Memory
UC15-HK	Peripheral Processor: 11/05 or 11/10 - NC or - SA, 2 DR11-C, DR15-C, MX15-B, DD11-B, KY11-JH, H950, MM11-K, 115V.	12K Local Memory
UC15-HL	Peripheral Processor: 11/05 or 11/10-ND or - SA, 2 DR11-C, DR15-C, MX15-B, DD11-B, KY11-JH, H950, MM11-K, 230V.	12K Local Memory
RK15-HE	RK05-AA, RK11-E, UC15-HE, 115V, 60Hz	
RK15-HF	RK05-BB, RK11-E, UC15-HF, 230V, 50Hz.	
RK15-HH	RK05-AB, RK11-E, UC15-HF, 230V, 60Hz.	
RK15-HJ	RK05-BA, RK11-E, UC15-HE, 115V, 50Hz.	
RK15-HK	RK05-AA, RK11-E, UC15-HK, 115V, 60Hz.	
RK15-HL	RK05-BB, RK11-E, UC15-HL, 230V, 50Hz.	
RK15-HM	RK05-AB, RK11-E, UC15-HL, 230V, 60Hz.	
RK15-HN	RK05-BA, RK11-E, UC15-HK, 115V, 50Hz.	
15/76-DE	KP15, ME15-EA, LA30-CA, PC15, KE15, KW15, TC15, TU56, RK15-HE, 115V, 60Hz.	
15/76-DF	KP15, ME15-EB, LA30-CD, PC15-A, KE15, KW15, TC15, TU56, RK15-HF, 230V, 50Hz.	
15/76-DK	KP15, ME15-EA, LA30-CA, PC15, KE15, KW15, TC15, TU56, RK15-HK, 115V, 60Hz.	
15/76-DL	KP15, ME15-EB, LA30-CD, PC15-A, KE15, KW15, TC15, TU56, RK15-HL, 230V, 50Hz.	
15/76-ME	KP15, ME15-EA, LA30-CA, PC15, KE15, KW15, TC59-D, TU10, RK15-HE, 115V, 60Hz.	

15/76-MF KP15, ME15-EB, LA30-CD, PC15-A, KE15, KW15, TC59-D,
TU10, RK15-HF, 230V, 50Hz.

15/76-MK KP15, ME15-EA, LA30-CA, PC15, KE15, KW15, TC59-D,
TU10, RK15-HK, 115V, 60Hz.

15/76-ML KP15, ME15-EB, LA30-CD, PC15-A, KE15, KW15, TC59-D,
TU10, RK15-HL, 230V, 50Hz.

CHAPTER 1

INTRODUCTION

The Magnetic Tape Dump (MTDUMP) Program is a utility program of the PDP-15 ADVANCED Software System which provides users of industry-compatible magnetic tape with functions which are peculiar to this medium. In general, the program provides magnetic tape users with functions similar to those found in PATCH and DUMP. In addition, the program complements PIP with regard to magnetic tape functions; however, few functions which could be performed by PIP are duplicated.

The program MTDUMP is device dependent and accomplishes all magnetic tape I/O with .TRAN and MTAPE System Macro instructions; it cannot be used with other I/O devices.

1.1 FUNCTIONS

The following paragraphs briefly explain the basic functions of MTDUMP. A summary of commands is provided in Appendix A.

1.1.1 Dump File

One of the most common requirements of the magnetic tape user is the ability to examine portions of a tape. The Dump File facility in MTDUMP is intended to meet that need in a general and useful way. Simply stated, the Dump File is the repository of (1) images of command lines received from the keyboard and (2) groups of ASCII lines which represent, in readable form, the contents of the tape being examined in response to typed requests. The contents and format of the file, however, are subject to considerable variation and, in fact, the destination of the file may itself be changed during the run.

1.1.2 File Modification

This feature provides a convenient means for file updating or patching. Individual records may be accessed, allowing each word in the record to become available for examination and modification. Words and entire records may be inserted or deleted from the file and new files may thus be created.

1.1.3 File Transfer

This function, consisting of one instruction, permits copying magnetic tape on a record-for-record basis.

1.1.4 Directory Listing

These commands permit rapid listing and clearing of magnetic tape directories.

1.2 I/O DEVICES

The program accesses a maximum of three devices: the teleprinter, used for command string input and error reports; the magnetic tape transports (via MTA. or MTF.) for all input and output to all magnetic tape units; and an optional third device which is the destination device for what is termed the "Dump Output File". This file may contain records of commands typed to the program and any hard-copy response to these commands (normally record-by-record dumps). Dump Output may be directed to any device, including a magnetic tape. If magnetic tape is used for this purpose, however, the unit assigned may not also be manipulated by commands to MTDUMP. If no Dump Output file is desired, the teleprinter should be assigned as the Dump Output device.

1.3 ADDING MTDUMP TO THE USER SYSTEM

The program MTDUMP and its associated handlers (i.e., MTA, MTC, and MTF) are supplied to the user on the ADVANCED Monitor System Peripheral DECTape (DEC-15-SZZB-UC). Users who wish more convenient access to MTDUMP should relocate the program onto the system device using the utility program PIP.

MTDUMP may also be added to the system device as a System Program, using the facilities provided by the SGEN and PATCH utility programs. Refer to PDP-15 manuals DEC-15-YWZA-DN3 and -DN5 for the procedures needed to install MTDUMP onto the system device as a System Program.

If MTDUMP is relocated to the system device by either of the above means, its associated magnetic tape handlers must also be added to the system library (.LIBR BIN). This is accomplished using the utility program UPDATE. The use of UPDATE to insert the handlers MTA, MTC, and MTF is demonstrated in the following example:

UPDATE V8A

>US+ (ALT MODE)	Request Options U and S
>I MTA.,DTC.)	Insert routine MTA after routine DTE
>I MTC.)	Insert routine MTC next
>IMTF.)	Insert routine MTF next
>CLOSE)	Terminate UPDATE operations.

Refer to the Utility Programs manual DEC-15-YWZA-D for a complete description of UPDATE and its use.

CHAPTER 2

OPERATING PROCEDURE

2.1 DEVICE ASSIGNMENTS

MTDUMP is supplied as a relocatable program (MTDUMP BIN) and is loaded by the Linking Loader. Before loading, the user must make the following .DAT slot assignments:

.DAT slot -4	The device from which MTDUMP is to be loaded. If the program is on magnetic tape, MTA on .DAT slot 1 requires MTA on .DAT slot -4; MTF on .DAT slot 1 requires MTC on .DAT slot -4.
.DAT slot 1	MTA \emptyset or MTF \emptyset
.DAT slot 3	The Dump Output device, if required; or TTA if no Dump Output File is wanted.

2.2 PROGRAM STARTUP

After loading, the program types on the teleprinter:

MTDUMP Vnn	where: "Vnn" is the current version
BUFSIZ m	and "m" is the total number
>	(in decimal) of registers
	available for I/O buffers.

Each time the program is ready to accept a keyboard command, a right angle bracket (>) is typed.

At start (or restart) time, all magnetic tape units are automatically set to transfer in odd parity at 8 $\emptyset\emptyset$ BPI and at the channel count given by .SCOM+4, bit 6 (\emptyset means 7-channel, 1 means 9-channel). The user must issue a new FORMAT request (see paragraph 3.2.1) to effect transfer in another (non-standard) mode.

2.3 PROGRAM RESTART

To restart MTDUMP, type CTRL P, which causes the program to close the Dump Output File (if open) on .DAT slot 3. Then repeat program startup procedure.

NOTE

If the Dump Output has been directed to the teleprinter, CTRL P is acted upon only after completion of current line of output. To effect immediate termination, type CTRL P CTRL U.

COMMANDS

3.1 COMMAND STRING

MTDUMP accepts commands from the Teletype in the general format shown below. Formats for specific commands may vary significantly from this and are shown in the descriptions of the individual commands.

MTDUMP command formats are variations of the following:

$$c _ u_1, u_2, t)$$

where:

- c is the name of the function wanted.
- u_1 is a digit specifying the source unit for two-unit operations (e.g., COPY) or the one object unit for single-unit operations.
- u_2 is a digit specifying the destination unit for two-unit operations or is absent for single-unit operations.
- t specifies a condition (either count overflow or transport status) which, when encountered, causes termination of the function whose name is "c". "t" may be absent and, if not given, is assigned the implicit integer value 1. Explicit values of "t" may include:
 - a. An integer in absolute value less than $262,144_{10}$ and greater than zero
 - b. The character string "EOT" (END OF TAPE)
 - c. The character string "BOT" (BEGINNING OF TAPE)
 - d. The character string "EOF" (END OF FILE)

Parameters are separated from the command by a space ($_$) and from each other by commas. The command line is terminated by a carriage return ($)$).

Some commands require only a single argument, while others require all three.

Example:

$$\text{REWIND} _ 1)$$

Only the single object unit need be specified; further, the terminating

condition "BOT" is implicit in the command and need not be given. Copying an entire logical tape from Unit 1 to Unit 2, however, requires all three parameters.

Example:

```
COPY_1,2,EOT)
```

3.1.1 Terminating Conditions

As indicated above, the "t" specification in the command line may be either an integer or a character string or absent. If "t" is an integer, the value of the numeric string represents the number of physical records to be treated during the operation requested.

Example:

```
SPACE_1,80)
```

This command string means: Evaluate the string "80" according to the radix currently in effect, then space the tape on drive 1 forward until that many records have been passed over. If, in the example, tape 1 was at loadpoint and if the prevailing radix was decimal, then at the completion of the operation the read/write head would be positioned between the 80th and 81st physical records on tape.

Example:

```
COPY_1,2,80)
```

The above example causes a transfer of 80 physical records from drive 1 to drive 2, leaving the read/write head on each drive positioned immediately following the last record transferred.

If "t" is a non-numeric string (EOT, BOT, EOF), then the operation requested is deemed complete when one of the following conditions is observed:

- a. EOT Two consecutive EOF markers have been passed in either reverse or forward direction.
- b. BOT The loadpoint marker has been reached (but not passed) in the reverse direction.
- c. EOF A single EOF marker has been passed in either direction.

If "t" is the string "EOF" or "EOT", the position of the read/write head relative to the EOF marker causing termination depends upon the direction of tape motion when the condition is encountered.

Example:

```
BACKSPACE l, EOF )
```

The read/write head will be positioned just before the marker. The next record read in the forward direction will be the EOF marker just passed in backspacing.

If "t" is the string "BOT", the head is left positioned just after the loadpoint; the program will not backspace over BOT.

If "t" is absent from a command string in which it is required, then the value 1 is assumed. Thus the commands in the following example are equivalent.

Example:

```
SPACE l, l )  
SPACE l )
```

3.1.2 Command Abbreviations

Most commands in MTDUMP may be abbreviated to a single letter (the initial character). In the command descriptions which follow, legal abbreviations are shown immediately following the command and enclosed in parentheses.

Example:

```
REWIND (R) u, t )
```

3.2 SETUP COMMANDS

This is a group of commands which generally apply to most major functions of MTDUMP. These commands are usually given prior to the execution of a function (e.g., DUMP, COPY).

3.2.1 Set Non-Standard Tape Format

The initial setup for input and output tapes is odd parity at 800 BPI (the channel count is given by .SCOM+4, bit 6). The FORMAT command allows the user to change the parity, density, and/or channel count.

Usage:

```
FORMAT (F) u, pd )
```

where: "u" is the tape whose format is being set and "pdc" is a group of three single-character parity, density, and channel-count indicators, as follows:

p (parity) is "E" (even) or "O" (odd)
d (density) is "2" (200 BPI), "5" (556 BPI), or
"8" (800 BPI)
c (channel) is "9" (9-channel) or "7" (7-channel)

The three descriptors may appear in any order, and any may be absent, in which case the relevant status for the tape remains unchanged.

Example:

```
FORMAT┐2,E57,)  
      or  
FORMAT┐2,5E7,)  
      or  
FORMAT┐2,75E,)
```

All of the above examples set up tape unit 2 for even parity, 556 BPI, 7-channel operation.

Example:

```
FORMAT┐2,O,)  
      or  
F┐2,O,)
```

These commands change the parity of tape unit 2 without disturbing the current density or channel count.

NOTE

The only legal density for a 9-channel tape drive is 800 BPI. Requests for other densities will not be honored.

FORMAT commands are effective until MTDUMP is restarted via the CTRL P function.

3.2.2 Set Standard Tape Format

Standard System Format may be requested for any unit. A special case of the FORMAT command is employed to unconditionally reset tape format to odd parity, 800 BPI, and 7- or 9-channel (according to .SCOM+4, bit 6.)

Usage:

FORMAT(F) \lfloor u, D \rfloor

where: "u" is the unit whose format is to be set and the character "D" means "default".

3.2.3 Specify Global Radix

The program always treats certain numeric strings (e.g., unit specification) as octal. Others, however, may be specified as either octal or decimal by the NUMBER command. The following numeric groups are interpreted (on input) or printed as octal or decimal strings according to the argument given in the latest NUMBER request:

- a. The "t" specification in command lines (where applicable) when "t" is an integer. If the current radix is octal, then the command:

SPACE \lfloor 1, 20 \rfloor

causes the tape on unit 1 to be spaced forward 16₁₀ records.

- b. The word sequence numbers of dumped data.
- c. The word sequence numbers of EXAMINE requests. (See below.)
- d. The record-length argument of the SIZE request. (See below.)

The radix specified remains in effect until another NUMBER command is encountered or the program is restarted. The default radix is octal.

Usage:

NUMBER(N) \lfloor { OCTAL
DECIMAL } \rfloor

3.2.4 Specify Local Radix

The radix of a number string in a single command line may be specified by a one-character suffix, D for decimal, K for octal. Such specification overrides the current global radix, but is in effect only during the processing of the command line in which the suffix appears. Local radix control may be used following:

- a. The "t" specification in command lines (where applicable) when "t" is an integer.
- b. The word sequence numbers of EXAMINE requests.

c. The record-length argument of the SIZE request.

Example:

```
SPACE└1,20D)
```

The command above causes tape unit 1 to space forward 20_{10} records regardless of the current global radix.

Example:

```
SPACE└1,20K)
```

Similarly, this command spaces the tape forward 20_8 (16_{10}) records.

3.2.5 Command-Line Echo

Legal keyboard requests are placed in the Dump Output File, exactly as typed, to allow the user to correlate the progress of the run, relative tape position, and the record contents during later examination of the hard-copy dump. Command-line echo can be bypassed, however, by use of the VERIFY command.

Usage:

```
VERIFY(V)└{ON  
          OFF}
```

If ON or OFF is not specified, ON is assumed.

Example:

```
v)
```

When MTDUMP is first loaded or is restarted, VERIFY mode is set ON. If the teleprinter is the assigned dump output device (.DAT slot 3), command-line echo is not performed. Illegal commands are not echoed.

3.2.6 Dump File Display Format

The input tape is output to the Dump File as individual physical records. Each record is represented as a number which indicates record length in ASCII lines. Each line, in turn, contains:

1. A sequence number which reflects the position in the record of the first data word in the line displayed.
2. A string of data words or data-word pairs.

Sequence numbers are in either octal or decimal notation; the radix is chosen in response to the last previous NUMBER command.

Display format is set by the MODE request followed by the appropriate argument.

Usage:

MODE $\left\{ \begin{array}{l} \text{OCTAL} \\ \text{SYMBOLIC} \\ \text{TRIMMED} \\ \text{ASCII} \end{array} \right\}$

Where:

OCTAL Displays single words as six octal digits.
 SYMBOLIC Displays single words as a three-character operation-code mnemonic, an "indirection" indicator (*), if present, and a 13-bit (5-digit) address.
 TRIMMED Displays single words as three six-bit alphanumeric characters.
 ASCII Displays pairs of words as five seven-bit ASCII characters. A blank is printed for each character outside the range $40_8 - 137_8$.

The default assumption is OCTAL and implicit in the request:

MODE,

The table below shows examples of data-word treatment in each of the four modes.

<u>OCTAL</u>	<u>SYMBOLIC</u>	<u>TRIMMED</u>	<u>ASCII</u>
512132 744634	AND 12132 OPR 04634)QZ <&\	REWIND
420320 000000	XCT*00320 CAL 00000	#CP 000	D
777777 010203	LAW 17777 CAL 10203	??? ABC	A

3.2.7 Inserting Comments in the Dump File

Explanatory notes may be placed in the output file by use of the LOG command. When the LOG request is encountered, subsequent typed input is taken as commentary and is added, exactly as it appears, to the Dump Output File. Carriage returns may be included, and multiple lines may be inserted with a single LOG request. An ALTMODE terminates each comment and causes the program to accept a new request.

Usage:

```
LOG└─comments)
comments.....)
(ALTMODE)
```

3.2.8 Return Control to Monitor

An EXIT request causes the program to close the Dump Output File (if one is open) on .DAT slot 3, then perform an .EXIT return to the Monitor. Use this command for return to the Monitor if the program is being run in the Batch Environment.

Usage:

```
EXIT)
```

3.3 MANIPULATIVE FUNCTIONS

The following commands position the tape and write EOF markers on the tape drive specified.

3.3.1 Rewind Tape

This command initiates a rewind on tape unit "u".

Usage:

```
REWIND(R)└─u)
```

3.3.2 Backspace Tape

This command backspaces the tape on unit "u" until the "t" condition is satisfied.

Usage:

```
BACKSPACE(B)└─u,t)
```

where: "t" is an integer (number of records), "EOF", "EOT", or "BOT".

3.3.3 Space Tape

This command spaces the tape on unit "u" forward until the "t" condition is satisfied.

Usage:

```
SPACE(S)└─u,t)
```

where: "t" is an integer (number of records), "EOF", or "EOT".

3.3.4 Write End-of-File Marker

This command writes a single "EOF" marker on tape unit "u".

Usage:

```
TAPEMARK(T)␣u,
```

3.4 DUMP FILE OPERATIONS

3.4.1 Dump File Management

The Dump Output File may be written on any physical device. If the device chosen is file-structured, however, the user must specify a name to be given the Dump File and must explicitly request that the file be closed (unless the EXIT command is used). Furthermore, the file name must be given before any other requests are issued.

Usage:

```
OPEN␣filename␣ext,
```

where: filename is the name of the file to be created.

ext is the filename extension. If omitted, "LST" is the default assumption.

If an OPEN request is not given, the program types

```
NO DUMP FILE OPEN  
>
```

on the Teletype and waits for another command.

NOTE

The comment is actually printed when an attempt is made to write into the Dump File, i.e., at command-line echo if VERIFY is ON or at Dump-Record Output if VERIFY is OFF.

A check is made to ensure that the filename given is unique. If a file of the name specified already exists on the Dump Output device, the program types:

```
FILE FOUND ON DUMP DEVICE:      filename ext  
DO YOU WISH TO DELETE IT?  
>
```

The program then waits for the user to type a response to the query.
Typing

)
or
Y)
or
YES,

indicates the affirmative, and the already-existing file is overlaid (i.e., deleted when the new file is .CLOSEd). Any other response is negative and the program returns to accept a new keyboard command.

The Dump Output File is closed upon receipt of the CLOSE command from the keyboard.

Usage:

CLOSE,

or whenever the program is restarted (CTRL P).

3.4.2 Dump Tape Records

This command dumps records from unit "u" into the named file open on .DAT slot 3. The sequencing of data words and the format in which they are written are controlled by the latest NUMBER and MODE requests.

Usage:

DUMP(D)␣u,t,

where: "u" is the tape unit number
"t" is an integer (number of records), "EOF", or "EOT".

3.4.3 Dump Tape Records on the Teleprinter

This command performs the same function as the DUMP command, except that the records are unconditionally dumped on the teleprinter.

Usage:

LIST(L)␣u,t,

3.4.4 Tape Status

In addition to data input from magnetic tape and the Teletype, the Dump Output File contains indicators of status encountered on the tape being read. Comments are added to the file (and typed on the teleprinter) in response to the following observed conditions on the tape.

<u>Message</u>	<u>Meaning</u>
*END OF FILE ENCOUNTERED	An unexpected end-of-file marker was read.
*PHYSICAL EOT ENCOUNTERED	The end-of-tape reflective spot was reached on input or output.
*BUFFER OVERFLOW	The tape record read is too long to be accommodated in the available buffer space.
*BOT ENCOUNTERED	The loadpoint reflective spot was unexpectedly reached during a backspace operation.
*PERMANENT READ ERROR ENCOUNTERED	After 64 ₁₀ read attempts, the input record still has not been transferred correctly. The read/write head is positioned immediately before the record.

3.4.5 Example of Dump Operation

The following example shows the instructions required to dump the file directory of magnetic tape unit Ø in octal format (to allow the accessibility map to be examined) and then in trimmed ASCII format (to allow reading of the file name entries).

Examples:

REWIND 0
SPACE 0,1
MODE OCTAL
DUMP 0,1

1	747377	000000	747750	777777	750000	000000	000000	000000	000000	000000	000000	131571
10	000000	000000	000000	000000	000000	000000	000000	000000	000000	000000	000000	000000
19	000000	233123	231320	021413	233123	111702	141300	233123	562331			
28	231404	233123	561411	022200	021116	561417	010400	021116	040424			
37	710000	021116	031001	111600	021116	302205	060000	021116	050105			
46	061114	021116	161716	061114	021116	252004	012405	233123	053005			
55	032524	233123	050411	240000	233123	201120	000000	233123	066400			
64	000000	233123	150103	221700	233123	066401	000000	233123	150103			
73	221701	233123	152423	070516	233123	031716	260000	233123	152402			
82	171724	232203	152404	251520	021116	000000	000000	000000	000000			
91	000000	000000	000000	000000	000000	000000	000000	000000	000000			
100	000000	000000	000000	000000	000000	000000	000000	000000	000000			
109	000000	000000	000000	000000	000000	000000	000000	000000	000000			
118	000000	000000	000000	000000	000000	000000	000000	000000	000000			
127	000000	000000	000000	000000	000000	000000	000000	000000	000000			
136	000000	000000	000000	000000	000000	000000	000000	000000	000000			
145	000000	000000	000000	000000	000000	000000	000000	000000	000000			
154	000000	000000	000000	000000	000000	000000	000000	000000	000000			
163	000000	000000	000000	000000	000000	000000	000000	000000	000000			
172	000000	000000	000000	000000	000000	000000	000000	000000	000000			
181	000000	000000	000000	000000	000000	000000	000000	000000	000000			
190	000000	000000	000000	000000	000000	000000	000000	000000	000000			
199	000000	000000	000000	000000	000000	000000	000000	000000	000000			
208	000000	000000	000000	000000	000000	000000	000000	000000	000000			
217	000000	000000	000000	000000	000000	000000	000000	000000	000000			
226	000000	000000	000000	000000	000000	000000	000000	000000	000000			
235	000000	000000	000000	000000	000000	000000	000000	000000	000000			
244	000000	000000	000000	000000	000000	000000	000000	000000	000000			
253	000000	000000	000000	000000	000000	000000	000000	000000	000000			

BACKSPACE 0,1
MODE TRIMMED
DUMP 0,1

1	< ; ?	@ @ @	< ? (? ? ?	= @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @
17	@ @ @	K M 9	@ @ @	S Y S	S K P	B L K	S Y S	I O B	L K @	S Y S	. S Y	S L D	S Y S	. L I	B R @	B I N	
33	. L O	A D @	B I N	O D T	9 @ @	B I N	C H A	I N @	B I N	X R E	F @ @	B I N	E A E	F I L	B I N	N O N	
49	F I L	B I N	U P D	A T E	S Y S	E X E	C U T	S Y S	E D I	T @ @	S Y S	P I P	@ @ @	S Y S	F 4 @	@ @ @	
65	S Y S	M A C	R O @	S Y S	F 4 A	@ @ @	S Y S	M A C	R O A	S Y S	M T S	G E N	S Y S	C O N	V @ @	S Y S	
81	M T B	O O T	S R C	M T D	U M P	B I N	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @
97	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @
113	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @
129	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @
145	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @
161	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @
177	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @
193	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @
209	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @
225	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @
241	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @	@ @ @
257	@ @ @																

3.5 TRANSFER FUNCTION

The COPY command allows the user to perform record-for-record copying of tapes.

Usage:

```
COPY(C) u1, u2, t
```

where: "u₁" is the source drive

"u₂" is the destination drive

"t" may be an integer (number of records), "EOF", or "EOT"

Standard parity and density (odd parity, 800 BPI) prevail, unless they have been changed by a FORMAT request.

To copy an entire tape from unit 1 to unit 2, for example:

```
REWIND u1)  
REWIND u2)  
COPY u1, u2, EOT)
```

To replace the last data record on unit 2 with the first data record on unit 1:

```
REWIND u1)           /find first record on 1.  
SPACE u2, EOT)       /find last record on 2.  
BACKSPACE u2, 3)     /backspace over two EOF's plus one data record.  
COPY u1, u2, 1)      /copy 1 record from 1 to 2.  
TAPEMARK u2)         /make a new EOT  
TAPEMARK u2)         / indicator on 2.
```

3.6 FILE MODIFICATION

The file modification feature of MTDUMP allows the user to access single records, modify or delete words in a record, delete entire records, or add new records to his file.

3.6.1 Read a Single Record

The next sequential physical record is read from tape unit "u" and is stored in core. Its length is saved in anticipation of a subsequent PUT request (see Paragraph 3.6.4).

Usage:

```
GET (G) u
```

At the completion of input, the following message is printed indicating

the length, in words, of the record just read.

RECSIZE:nn

3.6.2 Examine and Modify Data Words

Designed for use in conjunction with the GET and PUT commands, the EXAMINE request allows the user to access and update individual data words in the program buffer. Any number of contiguous registers may be examined and modified with a single command.

Usage:

EXAMINE(E)ln,

where: "n" is the relative position in the buffer (record) of the first word to be displayed. If a "D" or "K" suffix (see section 3.2.4) is present, the argument is interpreted appropriately. If no suffix is present, "n" is interpreted according to the current global radix. The argument specifies the position of a word relative to word \emptyset in the buffer.

Example:

EXAMINEl1,

The above command accesses the first data word in the buffer. The program responds to the command by displaying on the teleprinter the contents of the register specified in the mode (octal, symbolic, trimmed, ASCII) currently in effect. No carriage return is executed, however, after the displayed data word typeout. The user has several options.

- a. If a carriage-return is typed, the program responds by displaying the contents of the next higher register.
- b. If an ALTMODE is typed, buffer examination is deemed complete and the program returns to read a new command.
- c. A six-digit numeric string (octal notation) may be typed to replace the contents of the register being examined. The terminator of the line typed by the user may be either a carriage return or an ALTMODE. The terminator directs the program's activity after the desired modification has been performed. A carriage return opens the next sequential register; an ALTMODE returns control to the command processor.

3.6.3 Specify Output Record Length

The SIZE command specifies, in words, the length of the record to be written in response to a subsequent PUT request (see paragraph 3.6.4).

Usage:

```
SIZE [ n )
```

where: the parameter "n" is the total words in the output record. If a suffix "D" or "K" (see section 3.2.4) is present, the argument is evaluated appropriately. If no suffix is present, the numeric string is interpreted in the current global radix.

Output record size is implicitly set during input "GET" processing. The SIZE facility offers a means of overriding the implicit setting.

3.6.4 Write Single Record

The PUT command writes data residing in the program's buffer as the next sequential record on tape unit "u". The length of the record written is either the length of the record read in response to the latest GET request or the length specified in a SIZE request which occurred after the latest GET request.

Usage:

```
PUT(P) [ u )
```

3.7 DIRECTORY LISTING

This group of commands is available for dealing with the Magnetic Tape File Directory. The contents of the Directory on unit "u" may be printed on the teleprinter or written into the Dump Output File; and the Directory may be cleared. None of these commands may be abbreviated.

3.7.1 Write File Directory in Dump Output File

The contents of the File Directory of the tape specified are written in the Dump Output File.

Usage:

```
DDUMP [ u )
```

3.7.2 Print File Directory on Teleprinter

The contents of the File Directory of the tape specified are printed on the Teletype.

Usage:

DLIST u)

3.7.3 Clear Tape File Directory

Write a new (empty) File Directory on the tape specified.

Usage:

NEWDIR u)

APPENDIX A
SUMMARY OF COMMANDS

<u>COMMAND</u>	<u>MEANING</u>	<u>PARAGRAPH #</u>
<u>SETUP COMMANDS</u>		
EXIT,	Return Control to Monitor	3.2.8
FORMAT (F) <u>u</u> , pdc,	Set Non-Standard Tape	3.2.1
FORMAT (F) <u>u</u> , d	Set Standard Tape	3.2.2
LOG <u>comments</u> , comments....., (ALTMODE)	Insert one or more lines of comments	3.2.7
MODE { OCTAL SYMBOLIC TRIMMED ASCII }	Dump File Display	3.2.6
NUMBER { OCTAL DECIMAL }	Specify Global Radix	3.2.3

NOTE

D (decimal) or K (octal) specifies Local Radix which overrides Global Radix during processing of a single command line.

VERIFY (V) <u>ON</u> { ON OFF }	Bypass Command-Line	3.2.5
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MANIPULATIVE COMMANDS

BACKSPACE (B) <u>u</u> , t,	Backspace Tape	3.3.2
REWIND (R) <u>u</u> ,	Rewind Tape	3.3.1
SPACE (S) <u>u</u> , t,	Space Tape	3.3.3
TAPEMARK (T) <u>u</u> ,	Write End-of-File Marker	3.3.4

DUMP FILE COMMANDS

CLOSE,	Close Dump Output File	3.4.1
DUMP (D) <u>u</u> , t,	Dump Records into Named File	3.4.2
LIST (L) <u>u</u> , t,	Dump Records onto teleprinter	3.4.3
OPEN <u>filename</u> , <u>ext</u> ,	Open Named File	3.4.1

APPENDIX A (Cont.)

<u>COMMAND</u>	<u>MEANING</u>	<u>PARAGRAPH #</u>
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TRANSFER COMMAND

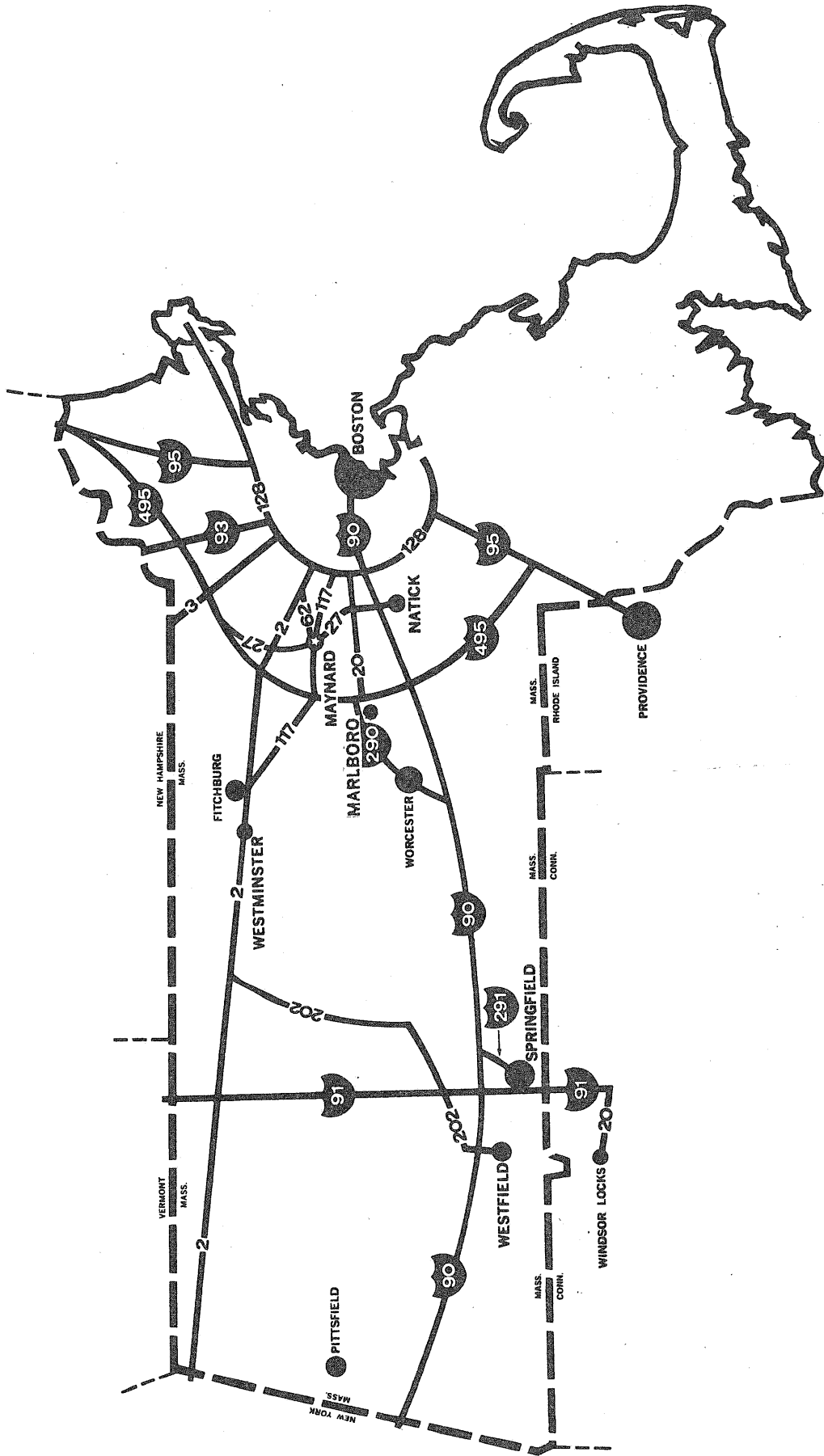
COPY(C)	u, u ₂ , t	Copy Tape Specified	3.5
---------	-----------------------	---------------------	-----

FILE MODIFICATION COMMANDS

EXAMINE(E)	n	Examine and Modify Data Words	3.6.2
GET(G)	u	Read Single Record	3.6.1
PUT(P)	u	Write Single Record	3.6.4
SIZE	n	Specify Output Record Length	3.6.3

DIRECTORY COMMANDS

DDUMP	u	Write File Directory in Dump Output File	3.7.1
DLIST	u	Print File Directory on Tele- printer	3.7.2
NEWDIR	u	Clear Tape File Directory	3.7.3



RTE. 20

EXIT 12

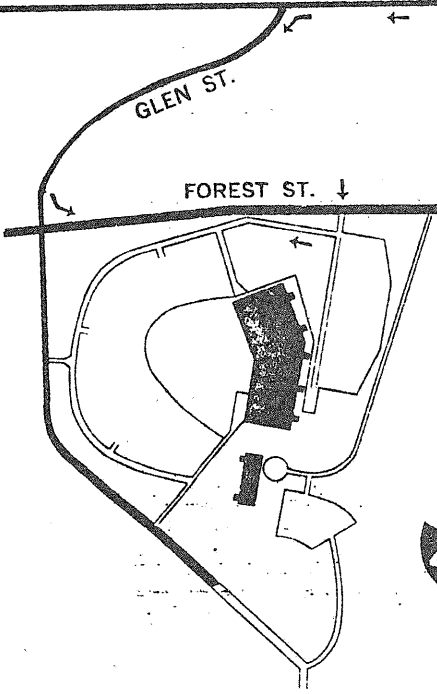
RTE. 20
TO WAYLAND

marlboro

GLEN ST.

FOREST ST. ↓

DIGITAL
MARLBORO
PLANT



495



EXIT 11

ROUTE 9

EXIT 11A

MASSACHUSETTS TURNPIKE



marlboro map

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The Software Communications Group, located at corporate headquarters in Maynard, publishes newsletters and Software Performance Summaries (SPS) for the various Digital products. Newsletters are published monthly, and contain announcements of new and revised software, programming notes, software problems and solutions, and documentation corrections. Software Performance Summaries are a collection of existing problems and solutions for a given software system, and are published periodically. For information on the distribution of these documents and how to get on the software newsletter mailing list, write to:

Software Communications
P. O. Box F
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Questions or problems relating to Digital's software should be reported to a Software Support Specialist. A specialist is located in each Digital Sales Office in the United States. In Europe, software problem reporting centers are in the following cities.

Reading, England	Milan, Italy
Paris, France	Solna, Sweden
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Tel Aviv, Israel	Munich, West Germany

Software Problem Report (SPR) forms are available from the specialists or from the Software Distribution Centers cited below.

PROGRAMS AND MANUALS

Software and manuals should be ordered by title and order number. In the United States, send orders to the nearest distribution center.

Digital Equipment Corporation Software Distribution Center 146 Main Street Maynard, Massachusetts 01754	Digital Equipment Corporation Software Distribution Center 1400 Terra Bella Mountain View, California 94043
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Outside of the United States, orders should be directed to the nearest Digital Field Sales Office or representative.

USERS SOCIETY

DECUS, Digital Equipment Computer Users Society, maintains a user exchange center for user-written programs and technical application information. A catalog of existing programs is available. The society publishes a periodical, DECUSCOPE, and holds technical seminars in the United States, Canada, Europe, and Australia. For information on the society and membership application forms, write to:

DECUS Digital Equipment Corporation 146 Main Street Maynard, Massachusetts 01754	DECUS Digital Equipment, S.A. P.O. Box 340 1211 Geneva 26 Switzerland
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Educational Services Department
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THE FOLLOWING PAGES ARE MISSING:

10-12

14-17

21-25

112-117

179-180

203-204

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