

# ***An Implementation of FFT, DCT, and Other Transforms on the TMS320C30***

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## Abstract

This book describes the several types of transforms and related algorithms used on the TMS320C30 family of digital signal processors. These include:

- The Fast Fourier Transforms (FFTs)
  - the complex radix-2 FFT
  - the complex radix-4 FFT
  - the real valued radix-2
- The Discrete Hartley Transform (DHT)
- The Discrete Cosine Transform (DCT)

The book contains:

- A description of transforms and their implementation on the TMS320C30 family of digital signal processors.
- A description and comparison of the different kinds of transforms: the FFTs, the Hartley transform and the Cosine transform
- A description of the features of the TMS320C30 that allow the efficient implementation of these algorithms
- Outlines of specific descriptions of implementations, transforms and TMS320C30 C Compiler facts
- Implementation issues
- Several graphics and tables detailing
  - Forms and flowgraphs of FFTs



- Memory requirements for FFT and Hartley transforms
- Differences in FFT and DCT timing

The end of the book contains 17 appendices with actual TMS320C30 source code for performing transforms.



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This report describes the implementation of several Fast Fourier Transforms (FFTs) and related algorithms on the TMS320C30. The TMS320C30 is the first device in the third generation of 32-bit floating-point Digital Signal Processors (DSPs) in the Texas Instruments TMS320 family. The algorithms considered here are the complex radix-2 FFT, the complex radix-4 FFT, the real-valued radix-2 FFT (both forward and inverse transforms), the Discrete Hartley Transform (DHT), and the Discrete Cosine Transform (DCT). These transforms have many applications, such as in image processing, sonar, and radar.

The introduction briefly describes transforms and their implementation on the TMS320 family of processors. Next, the different kinds of FFTs (including the real FFT), the closely-related Hartley transform, and the Cosine transform are described and compared. This is followed by a description of the TMS320C30 features that permit efficient implementations of these algorithms. Then, specific implementations, transforms, and TMS320C30 C Compiler facts are outlined. Finally, the report discusses some implementation issues, and the appendices list actual TMS320C30 code for performing transforms.

The powerful architecture and instruction set of the TMS320C30 permit flexible and compact coding of the algorithms in assembly language while preserving close correspondence to a high-level language implementation. The efficiency of the architecture and the speed of the device make faster realization of real and complex transforms possible. With the availability of a C compiler, these routines can be put in C-callable form and used as faster versions of FFT C functions.

## **Introduction**

The Fast Fourier Transform (FFT) is an important tool used in Digital Signal Processing (DSP) applications. Its development by Cooley and Tuckey gave impetus to the establishment of DSP as an independent discipline. The well-structured form of the FFT has also made it one of the benchmarks in assessing the performance of number-crunching devices and systems.

In recent years, because of the popularity of this signal-processing tool, there have been efforts to improve its performance by advances both at the algorithmic level and in hardware implementation. Researchers have been developing efficient algorithms to increase the execution speed of FFTs while keeping requirements for memory size low. On the other hand, developers of VLSI systems are including features in their designs that improve system performance for applications requiring FFTs. In particular, single-chip programmable DSP devices, currently available or under development, can realize FFTs with speeds that allow the implementation of very complex systems in realtime.



The Texas Instruments TMS320 family consists of five generations of programmable digital signal processors. The TMS32010 introduced the first generation, which today encompasses more than twelve devices with various speeds, interfacing capabilities, and price/performance combinations. FFT implementations on the TMS32010 can be found in the appendix of the book by Burrus and Parks [1].

The second-generation TMS320 devices (the TMS32020, the TMS320C25, and their spinoffs) enhanced the architecture and speed capabilities of the first generation. Examples of FFT programs implemented on the TMS32020 can be found in an application report in the book *Digital Signal Processing Applications with the TMS320 Family* [2]. Such programs are easily extended to the TMS320C25 because of the code compatibility between devices.

The architectural and speed improvements on the processors from one generation to the next have made the FFT computation faster and the programming easier. These advantages have reached a new high level in the third generation. The TMS320C30 is the first device in the third generation, and this report examines implementation of the FFT algorithms on it. The fourth generation (TMS320C4x) is a new set of floating-point devices, while the fifth generation (TMS320C5x) is a continuation of the fixed-point devices. Since software compatibility is maintained within the fixed-point and the floating-point devices, the existing FFT implementations will also be applicable to these new generations.

The Fourier Transform of an analog signal  $x(t)$ , given as

$$X(\omega) = \int_{-\infty}^{\infty} x(t) e^{-j\omega t} dt \quad (1)$$

determines the frequency content of the signal  $x(t)$ . In other words, for every frequency, the Fourier transform  $X(\omega)$  determines the contribution of a sinusoid of that frequency in the composition of the signal  $x(t)$ . For computations on a digital computer, the signal  $x(t)$  is sampled at discrete-time instants. If the input signal is digitized, a sequence of numbers  $x(n)$  is available instead of the continuous-time signal  $x(t)$ . Then, the Fourier transform takes the form

$$X(e^{j\omega}) = \sum_{n=-\infty}^{\infty} x(n) e^{-j\omega n} \quad (2)$$

The resulting transform  $X(e^{j\omega})$  is a periodic function of  $\omega$ , and it needs to be computed for only one period. The actual computation of the Fourier transform of a stream of data presents difficulties because  $X(e^{j\omega})$  is a continuous function in  $\omega$ . Since the transform must be computed at discrete points, the properties of the Fourier transform led to the definition of the *Discrete Fourier Transform* (DFT), given by

$$X(k) = \sum_{n=0}^{N-1} x(n) e^{-\frac{j2\pi kn}{N}} \quad (3)$$

When  $x(n)$  consists of  $N$  points  $x(0), x(1), \dots, x(N-1)$ , the frequency-domain representation is given by the set of  $N$  points  $X(k), k=0, 1, \dots, N-1$ . Equation (3) is often written in the form

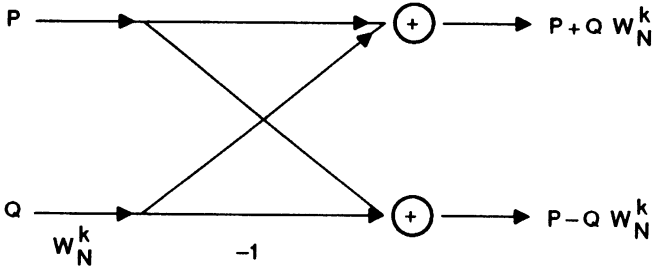
$$X(k) = \sum_{n=0}^{N-1} x(n) W_N^{nk} \quad (4)$$

where  $W_N^{nk} = e^{-j2\pi nk/N}$ . The factor  $W_N$  is sometimes referred to as the *twiddle factor*. A detailed description of the DFT can be found in references [1,3,4]. The computational requirements of the DFT increase rapidly with increasing block size  $N$ , having an impact on the real-time system performance. This problem was alleviated with the development of special fast algorithms, collectively known as Fast Fourier Transform (FFT). With an FFT, the computational burden increases much less rapidly with  $N$ , and for any given  $N$ , the FFT computational load, measured in terms of required multiplications and additions, is smaller than a brute-force computation of the DFT.

The definition of the FFT is identical to the DFT: only the method of computation differs. To achieve the efficiency of an FFT, it is important that  $N$  be a highly composite number. Typically, the length  $N$  of the FFT is a power of 2:  $N = 2^M$ , and the whole algorithm breaks down into a repeated application of an elementary transform known as a *butterfly*. If  $N$  is not a power of 2, the sequence  $x(n)$  is appended with enough zeroes to make the total length a power of 2. Again, references [1,3,4] contain a detailed development of the FFT. Reference [2] also discusses the same topic.

## Different Forms of the FFT

Over the years, researchers have developed different forms of FFT for more efficient computation. Special cases, such as those in which the input is a sequence of real numbers, have been investigated, and even more sophisticated algorithms have been developed. The general form of the FFT *butterfly* is given in Figure 1.



**Figure 1. Radix-2 Butterfly for Decimation in Time**

If the inputs to the butterfly are the two complex numbers  $P$  and  $Q$ , the outputs will be the complex numbers  $P'$  and  $Q'$ , such that

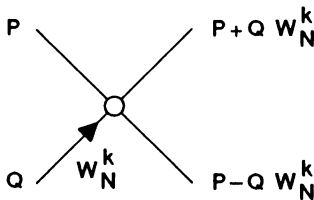
$$P' = P + Q W_N^k \tag{5}$$

and

$$Q' = P - Q W_N^k \tag{6}$$

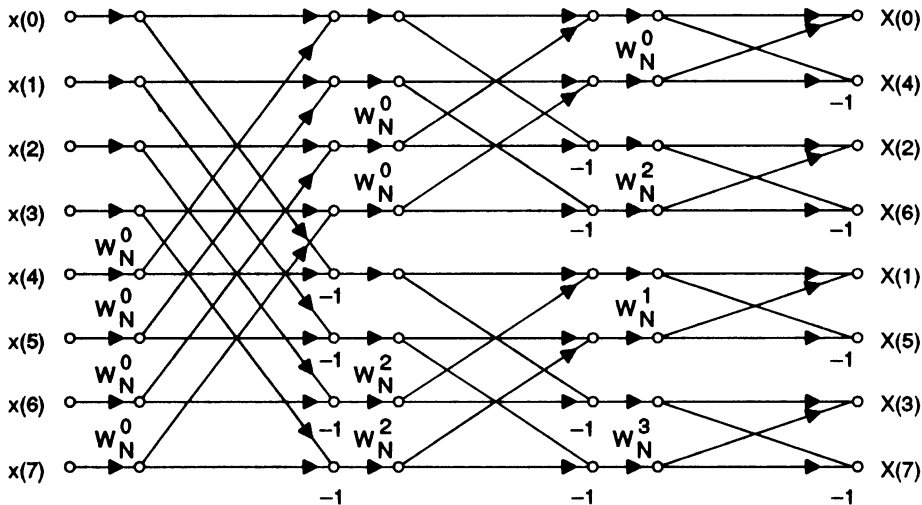
The quantities  $P$ ,  $Q$ , and  $P'$ ,  $Q'$  represent different points in the array being transformed, and they may or may not occupy adjacent locations in that array. In an in-place computation, the result  $P'$  will overwrite  $P$ , and  $Q'$  will overwrite  $Q$ .  $W_N^k$  represents again the twiddle factor, and its exponent is determined by the location of the corresponding butterfly in the FFT algorithm.

Figure 2 shows an alternate form of the same FFT butterfly.



**Figure 2. Alternate Form of Radix-2 Butterfly for Decimation in Time.**

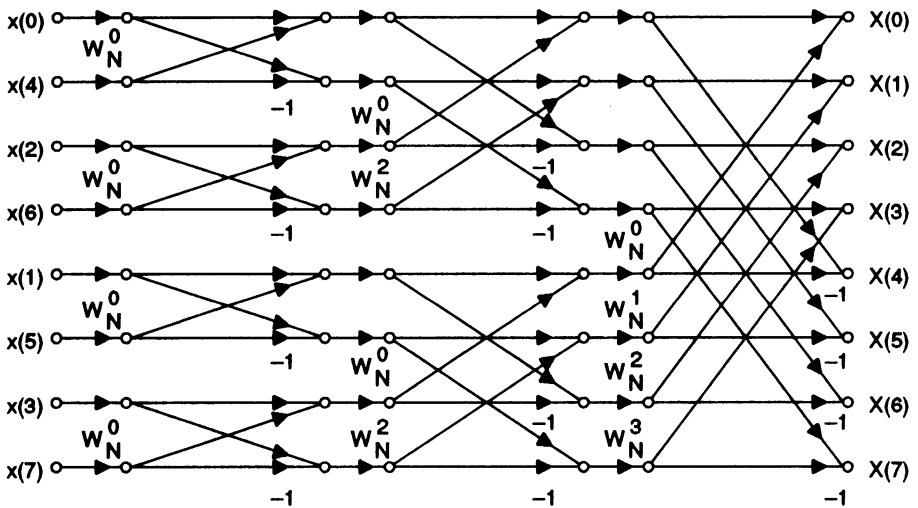
Although the notation is now less descriptive, it creates a clearer picture when several butterflies are put together to form an FFT. Using the first notation, Figure 3 is the flowgraph of an 8-point FFT example.



**Figure 3. Example of 8-Point FFT with Decimation in Time.**

Note that the input sequence  $x(n)$  is in the correct order, while the output  $X(k)$  is scrambled. Actually, this scrambling occurs in a very systematic way, called bit-reversed order: If you express the indices of a scrambled sequence in binary and you reverse this number, the result is the order that this particular point occupies. For instance,  $X(3)$  occupies the sixth position in the output (when counting from the zero position). In binary form,  $3_{10} = 011_2$ , and if bit-reversed, you get  $110_2 = 6_{10}$ , which is the position that  $X(3)$  occupies. It turns out that the third position is occupied by  $X(6)$ , and to restore the correct order at the output, you need only to swap these two numbers.

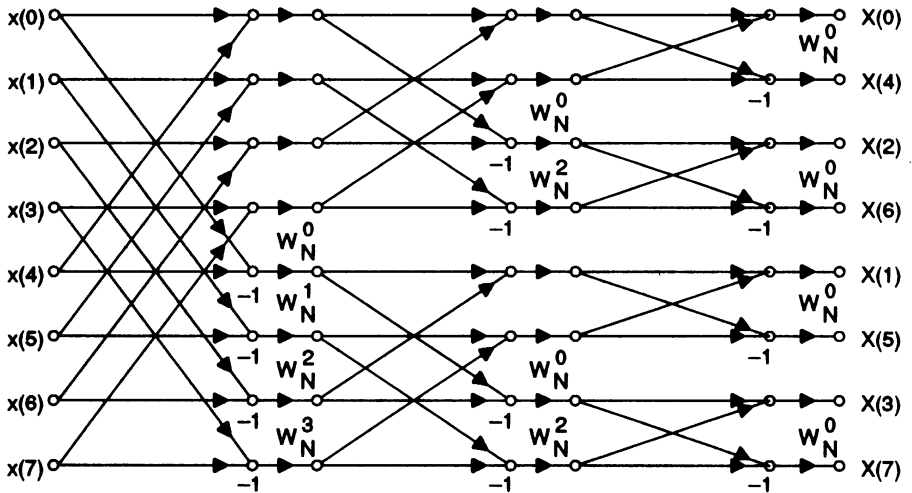
The same procedure can be repeated with all the scrambled numbers not occupying the position that their index suggests. If the input sequence  $x(n)$  is rearranged to appear in bit-reversed form, the output  $X(k)$  appears in the correct order, as shown in Figure 4.



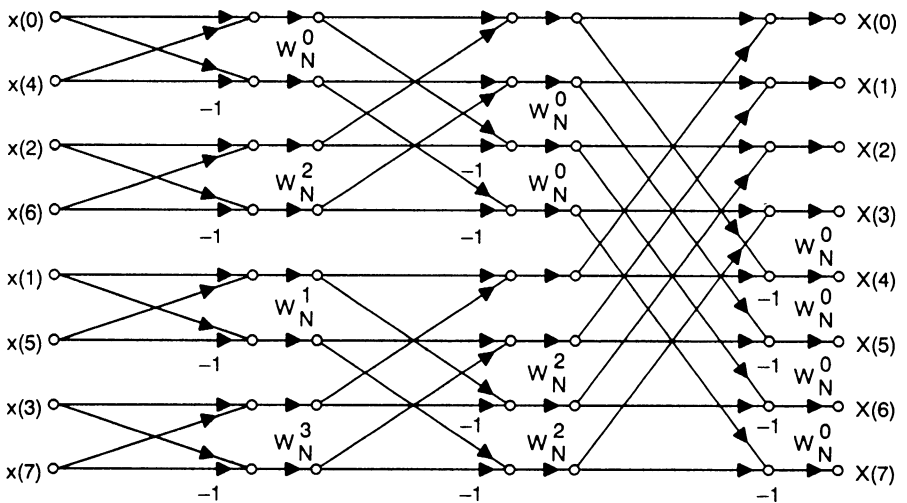
**Figure 4. Alternate Form of 8-Point FFT with Decimation in Time. The Input Is in Bit-Reversed Order and the Output Is in the Correct Order.**

Since the only difference between Figures 3 and 4 is a rearrangement of the butterflies, the computational load and the final results are identical. In terms of implementation, this rearrangement means that the nesting of the two innermost loops in the FFT routine is interchanged.

The butterflies and the FFT configurations presented thus far implement the FFT with a *decimation in time*. This terminology essentially describes a way of grouping the terms of the DFT definition; see Equation (3). An alternative way of grouping the DFT terms together is called *decimation in frequency*. Figures 5 and 6 show the same example of an 8-point FFT: Figure 5 with the input in correct order and the output in bit-reversed order, and Figure 6 vice-versa, and using the decimation in frequency (DIF).



**Figure 5. Example of an 8-Point FFT with Decimation in Frequency.**

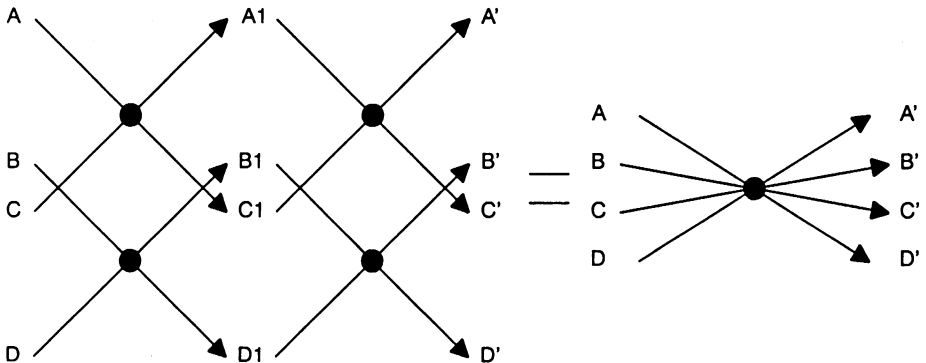


**Figure 6. Alternate Form of 8-Point FFT with Decimation in Frequency. The Input Is in Bit-Reversed Order and the Output Is in the Correct Order**

Pictorially, the difference between decimation in time and decimation in frequency is that the twiddle factor appears at the input of the butterfly in the first, and at the output in the second. Otherwise, the two methods are identical in terms of results. However, depending on what is the most convenient order of getting the twiddle factors and where the longest-span butterfly appears, you may prefer one method over the other.

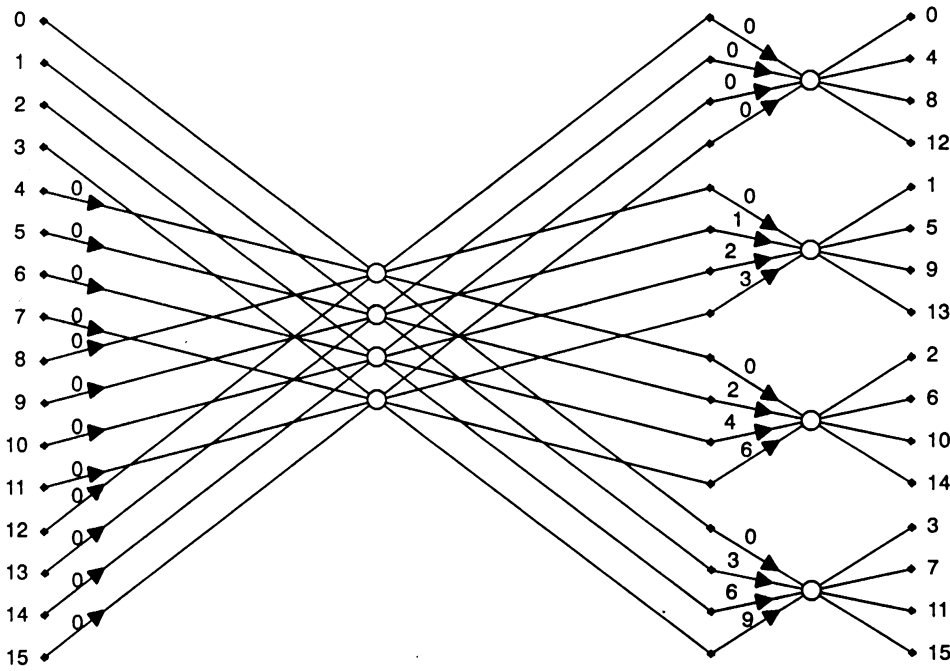
The butterfly shown in Figure 1 (or Figure 2) is the smallest element in a radix-2 FFT. The radix of the FFT represents the number of inputs that are combined in a butterfly. The Fast Fourier Transform is usually explained around the radix-2 algorithm for conceptual simplicity. If, however, higher-order radices are used, more computational savings can be achieved. These savings increase with the radix, but there is very little improvement above radix 4. That's why the radix-2 and radix-4 FFTs are the most commonly used algorithms.

In radix-4 FFT, each butterfly has 4 inputs and 4 outputs, essentially combining two stages of a radix-2 algorithm in one. Figure 7 shows this combination graphically.



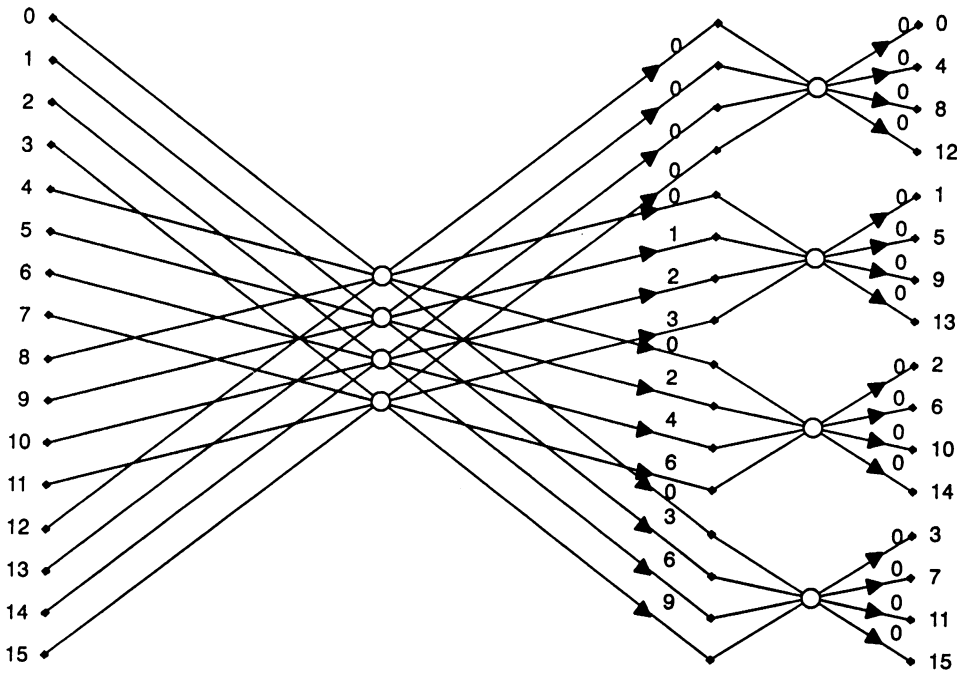
**Figure 7. Butterfly for Radix-4, Decimation-in-Time FFT.**

Although four radix-2 butterflies are combined into one radix-4 butterfly, the computational load of the latter is less than four times the load of a radix-2 butterfly. Examples of radix-4, 16-point FFTs are shown in Figures 8 and 9 for decimation in time and decimation in frequency, respectively.



**Figure 8. Example of a 16-Point, Radix-4, Decimation-in-Time FFT.**

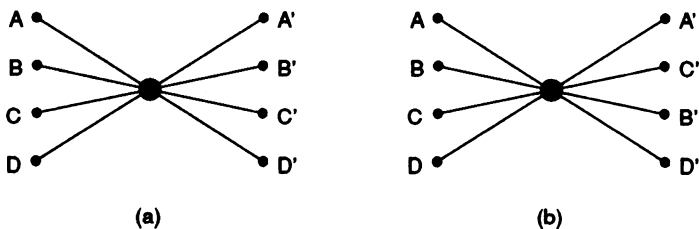




**Figure 9. Example of a 16-Point, Radix-4, Decimation-in-Frequency FFT.**

These configurations take the incoming sequence in order and produce the frequency-domain result in digit-reversed form. It is a simple matter to rearrange the FFT and have the input in digit-reversed form and the output in order.

*Digit reversal* is similar to bit reversal, except that the number whose digits are reversed is written in base 4 (equal to the radix) rather than base 2. For example, the output value  $X(14)$  in a 16-point, radix-4 FFT occupies position eleven (again starting from zero) because  $14_{10} = 32_4$  and, reversing the digits of the number,  $23_4 = 11_{10}$ . To restore the output to the correct order, the contents of locations with digit-reversed indices should be swapped. However, since the TMS320C30 has a special bit-reversed addressing mode, it is desirable to have the output of the radix-4 computation in bit-reversed rather than digit-reversed form. This is accomplished quite simply if, in each radix-4 butterfly, the two middle output legs are interchanged. That is, whenever the output of the butterfly is the four numbers  $A'$ ,  $B'$ ,  $C'$ , and  $D'$ , instead of storing them in that order, store them in the order  $A'$ ,  $C'$ ,  $B'$ , and  $D'$ , as shown in Figure 10.



**Figure 10. Radix-4 Butterflies. (a) Regularly-Ordered Output, (b) Bit-Reversed Output.**

References [5, 6] explain why this simple rearrangement puts the result in bit-reversed order.

### **Features of the TMS320C30**

The TMS320C30 is the first device introduced in the third generation of the TMS320 Digital Signal Processors [7,8]. It has many architectural features that permit very efficient implementation of algorithms. Some of those features pertinent to the FFT implementation are discussed in this section.

The two most salient characteristics of the TMS320C30 device are its high speed (60-ns cycle time) and floating-point arithmetic. The higher speed makes the implementation of real-time application easier than in earlier processors, even when the other architectural advantages are not considered. Each instruction executes in a single cycle under mild pipeline restrictions. The device automatically takes care of any potential conflicts. The pipeline should be observed closely (e.g., using the trace capability of the simulator) only if code optimization for speed is required.

The floating-point capability permits the handling of numbers of high dynamic range without concern for overflows. In FFT programs, in particular, the computed values tend to increase from one stage to the next, as discussed in reference [2]. Then, the fixed-point arithmetic will cause overflows if the incoming numbers are large enough and no provisions are made for scaling. All these considerations are eliminated with the floating-point capability of the TMS320C30. The TMS320C30 performs floating-point arithmetic with the same speed as any fixed point operation; no performance is sacrificed for this feature.

There are eight extended-precision registers, R0—R7, that can be used as accumulators or general-purpose registers, and eight auxiliary registers, AR0—AR7, for addressing and integer arithmetic. For many applications, these registers are sufficient for temporary storage of values, and there is no need to use memory locations. This is the case with the radix-2 FFT algorithm, where no locations are required other than those for the transformation of incoming data to be transformed. Also, arithmetic using these registers greatly increases the programming efficiency. The two index registers, IR0 and IR1, are used for indexing the contents of the auxiliary registers AR0—AR7, thus making the access of the butterfly legs and the twiddle factors easy.

A powerful structure in the TMS320C30 is the block-repeat capability that has the form

```
          RPTB    LABEL
          put instructions here
LABEL    last instruction
```

Whatever occurs after the RPTB instruction and up to the LABEL is repeated one time more than the number included in the repeat counter register, RC. The RC register must be initialized before entering the block-repeat construct. The net effect is that the repeated code behaves as if it were straight-line coded (no penalty for looping), with program size equal to the one in looped code. In this way, the FFT butterfly, being the core of the program, can be implemented in a block-repeat form, thereby saving execution time while preserving the clarity of the program and conserving program space.

A bit-reversed addressing mode is available to eliminate the need for swapping memory locations at the beginning or the end of the FFT (depending on the FFT type). When you use this addressing mode, you access a sequence of data points in bit-reversed order rather than sequentially, and you can recover the points in the correct order during retrieval of the data instead of spending extra cycles to accomplish it in software.

## **Implementation of Radix-2 and Radix-4 Complex FFTs**

Because of the powerful architecture and the instruction set of the TMS320C30, the assembly language program follows closely the flow of a high-level language program; this makes it easy to read and debug. It also keeps the size of the program small and reduces the requirements for program memory. Appendix A presents an example of code for a Radix-2 complex FFT, while Appendix B is a radix-4 complex FFT. The program memory requirements for these programs (as well as others to be discussed later) are given in Table 1.

**Table 1. Program Memory Requirements for the Core of the FFT and Hartley Transforms**

Routine Type	Program Size
Radix-2, complex FFT	50 words
Radix-4, complex FFT	170 words
Radix-2, real FFT	68 words
Radix-2, real inverse FFT	76 words
Hartley transform	71 words

The numbers in the table correspond only to the core program and do not include the sine/cosine tables for the twiddle factors, any input/output, or any bit-reversing operations. Note also that they are independent of the FFT data size.

The data memory requirements are, of course, dependent on the FFT size. The maximum length of a complex, radix-2 FFT that can be implemented entirely on the internal memory of the TMS320C30 is 1024 points. In the present implementation, the 1024-point radix-4 FFT requires a few more locations (about 7) than are available on-chip.

The code (provided in the appendices) has been written to be independent of the FFT length. The length  $N$ , together with the sine/cosine tables for the twiddle factors, should be provided separately to maintain the generic nature of the core FFT program. An example of a file with the sine/cosine tables for a 64-point FFT is given in Appendix F. Note that the FFT size and the number of stages are declared `global` in both files (i.e., the main routine and the file with the table) so that the core program gets the actual values during linking.

To reduce the storage requirements of a sine/cosine table, a full sine and a cosine cycle are overlapped. The table stores 5/4 of a full sine wave, with the cosine table starting with a phase delay of 1/4 cycle from the sine table. This table size is larger than actually needed, and it is selected merely for testing convenience of the algorithms. The minimum table size for a radix-2 complex FFT includes 1/2 of a full sine wave, and 1/2 of a full cosine wave. If these two half waves are combined using the above quarter-cycle phase delay, the minimum table size for this kind of FFT is 3/4 of a full sine wave. For instance, for a 1024-point FFT, the table can be the first 768 points of a sine wave, where a full cycle would be 1024 points. In the case of a radix-4 complex FFT, the minimum table size should include 3/4 of a sine and 3/4 of a cosine wave. Overlapping these requirements, we get the minimum table size of a radix-4 algorithm to be one full sine wave.

An example of a linking file is also included in Appendix F to show how the different segments are assigned. For a complete description of the assembler and linker, consult the corresponding manual [6].

The timing of the FFT routines was done using the cycle-counting capability of the TMS320C30 simulator. For the conversion of the number of cycles into seconds, a cycle time of 50 ns was used. The timing refers only to the core FFT computation, ignoring read-in and write-out requirements, since such requirements are application-dependent. Also, no bit reversal is counted (although it may be included in the program), since it is performed as part of the read-in or read-out. Table 2 gives the timing for the different FFT routines and for the Hartley transform.

**Table 2. FFT Timing in Milliseconds<sup>†</sup>**

Transform Size	Radix-2 Complex FFT	Radix-4 Complex FFT	Radix-2 Real FFT	Radix-2 Real Inverse FFT	Hartley Transform
64	0.101	0.103	0.047	0.053	0.068
128	0.211	—	0.099	0.110	0.151
256	0.453	0.520	0.215	0.241	0.336
512	0.991	—	0.476	0.535	0.943
1024	2.175	2.533	1.055	1.193	2.025
1024	1.972				

<sup>†</sup>Improvements have been made and are shown in this table. You may obtain the latest code from the BBS, (713) 274-2323.

The last entry in this table represents the timing of the radix-2, DIT routine generated at the University of Erlangen [18] and given in Appendix A. These numbers are typically used for benchmarking.

## Implementation of Real FFT

The development of FFT algorithms is centered mostly around the assumption that the input sequence consists of complex numbers (as does the output). This assumption guarantees the generality of the algorithm. However, in a large number of actual applications, the input is a sequence of real numbers. If this condition is taken into consideration, additional computational savings can be achieved because the FFT of a real sequence demonstrates the following symmetries: Assuming that the FFT output  $X(k)$  is complex,

$$X(k) = R(k) + j I(k) \quad (7)$$

and that the sequence has length  $N$ ,  $R(k)$  and  $I(k)$  should satisfy the following relations:

$$R(k) = R(N-k), \quad k = 1, \dots, N/2-1 \quad (8)$$

$$I(k) = -I(N-k), \quad k = 1, \dots, N/2-1 \quad (9)$$

$$I(0) = I(N/2) = 0. \quad (10)$$

In other words, the real part of the transform is symmetric around zero frequency, while the imaginary part is antisymmetric. Similar conditions hold if the transform is expressed in terms of magnitude and phase.

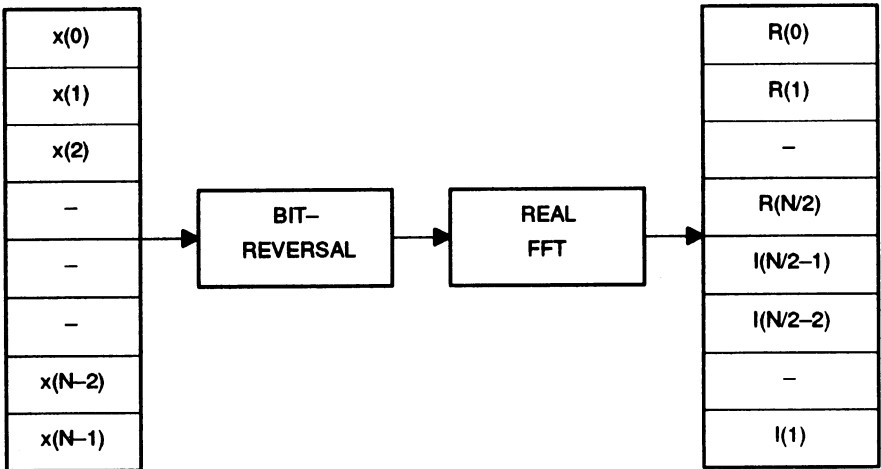
The savings are due to the fact that not all points need to be computed. Since the not-computed points do not need to be saved either, there are also storage savings. An efficient algorithm for real-valued FFTs is described in [10]. This algorithm was implemented in the present study in such a way that, given the sequence of  $N$  real numbers  $x(0), x(1), \dots, x(N-1)$ , the resulting FFT, consisting of complex numbers, is stored as  $R(0), R(1), \dots, R(N/2), I(N/2-1), I(N/2-2), \dots, I(1)$ .  $R(k)$  and  $I(k)$  represent the real and imaginary parts of the complex number  $X(k)$ . Figure 11 shows the memory arrangement for the FFT. Note that the input to the real FFT should be bit-reversed, but the bit reversal can be done as the data is brought in. With this arrangement, an  $N$ -point FFT uses exactly  $N$  memory locations. If the full array  $X(k)$  is needed, the following relations should be used:

$$X(0) = R(0) \tag{11}$$

$$X(k) = R(k) + j I(k), \quad K = 1, \dots, N/2-1 \tag{12}$$

$$X(N/2) = R(N/2) \tag{13}$$

$$X(k) = R(N-k) - j I(N-k), \quad k = N/2+1, \dots, N-1 \tag{14}$$



**Figure 11. Memory Arrangement of a Real FFT.**

It is expected that, in most signal processing applications, there will be no need to reconstruct the full  $X(k)$  array and that the output shown in Figure 11 will be sufficient for any further processing.

Appendix C contains TMS320C30 routines implementing a radix-2 real FFT and its inverse. The implementation of the forward transformation is based on the FORTRAN programs contained in [10]. The inverse transformation assumes that the input data are given in the order presented at the output of the forward transformation and produces a time signal in the proper order (i.e., bit-reversing takes place at the end of the program). Viewed another way, the inverse real FFT operates as shown in Figure 11 but with the arrows reversed (and inverse FFT taking the place of the FFT).

The timing for the real-valued FFT (both forward and inverse) is included in Table 2, and the corresponding program sizes are shown in Table 1. As you can see, the real-valued FFT is considerably faster than the corresponding complex FFT because not all the computations need be performed. Furthermore, there are data storage savings because only half the values must be stored. As a result, the maximum length of real-valued FFT that can be implemented on the TMS320C30 without using any external memory is 2048 points. Of course, if all the values are needed, they can be recovered using the symmetry conditions mentioned earlier. To achieve the efficiencies of real FFT and not use any extra memory locations during the computation, the decimation-in-time method is applied [10]. Decimation in time requires the bit-reversal operation in the forward transform to be performed at the beginning of the program rather than at the end. The reverse is true for bit-reversing in the inverse transform.

## The Discrete Hartley Transform

Another transform that has attracted attention recently is the Discrete Hartley Transform (DHT)[11, 12]. The DHT is applicable to real-valued signals and is closely related to the real-valued FFT. Comparison of references [10] and [12] describing the implementation of the two algorithms on FORTRAN programs shows that their implementation on the TMS320C30 should be similar. And indeed, this is the case.

The DHT pair is defined for a real-valued sequence  $x(n)$ ,  $n = 0, \dots, N-1$ , by the following equations:

$$H(k) = \sum_{n=0}^{N-1} x(n) \text{cas}(2\pi k n / N), \quad k=0, \dots, N-1 \quad (15)$$

$$x(n) = \frac{1}{N} \sum_{k=0}^{N-1} H(k) \text{cas}(2\pi k n / N), \quad k=0, \dots, N-1 \quad (16)$$

where  $\text{cas}(x) = \cos(x) + \sin(x)$ . The DHT demonstrates a symmetry that is convenient for implementations: The same program can be used for both the forward and the inverse transforms, and the result is correct within a scale factor. Also, the real FFT and the DHT can be derived from each other [12].

A radix-2 Hartley transform was implemented on the TMS320C30, and the corresponding code is included in Appendix D. This code follows the structure of the real FFT in Appendix C. Tables 1 and 2 show the program memory requirements and the timing for the execution of Hartley transforms of different sizes. The sine/cosine table sizes are the same as in the case of a real FFT.

## The Discrete Cosine Transform

The Discrete Cosine Transform (DCT), since its introduction in 1974 [13], has gained popularity in speech and image processing applications because of its near-optimal behavior. This discussion is based on the paper by Lee [14]. The DCT code was developed and implemented by Paul Wilhelm of the University of Washington.

If  $x(n)$ ,  $n=0, \dots, N-1$  is a time-domain signal and  $X(k)$  is the corresponding DCT,  $x(n)$  and  $X(k)$  are related by the following equations:

$$x(k) = \frac{2}{N} \sum_{n=0}^{N-1} e(k) x(n) \cos \frac{(2k+1)\pi n}{2N} \quad (17)$$

$$x(n) = \sum_{k=0}^{N-1} e(k) X(k) \cos \frac{(2k+1)\pi n}{2N} \quad (18)$$

$$e(0) = 1/\sqrt{2} \quad (19)$$

$$e(k) = 1, \quad \text{for } k \neq 0 \quad (20)$$

Appendix E shows an implementation of the DCT based on the paper by Lee [14]. The appendix contains the algorithms for both the forward and the inverse transformations and an example of a table for a 16-point DCT. Note that, because of the structure of the algorithm, the cosine table needed contains actually the inverses of the cosines (within a scale factor), and it is not stored in the natural order. Instead, it is generated by the following C pseudocode:

```
for [k=2, i=0; k=N/2; k*=2]
  for [j=k/2; j<N/2; j+=k]{
    cos__table[i++] = 1/[2*cos[j*pi/[2*N]]];
    cos__table[i++] = 1/[2*cos[(N-j)*pi/[2*N]]];
  }
cos__table[N-2] = cos[pi/4];
cos__table[N-1] = 2/N;
```



The last entry to the table is not part of the cosine itself; it is a constant that is used by the algorithm, and it is placed at the end of the cosine table for convenience.

Table 3 shows the timing of the forward and inverse transforms for different transform lengths. The difference in the timing between the forward and the inverse transforms is due to the fact that more time was expended to optimize the performance of the inverse transform. Since four of the smallest butterflies were done simultaneously in the center program loop, the minimum permissible array size to be transformed is 8.

**Table 3. DCT Timing in Milliseconds**

Transform Size	Forward Transform	Inverse Transform
16	0.019	0.017
64	0.875	0.073
128	0.192	0.161
256	0.418	0.347
512	0.912	0.754
1024	1.982	1.652

### Other Related Transforms

In addition to the FFT types mentioned earlier (complex, real, decimation-in-time, decimation-in-frequency, etc.), newer forms of the FFT have been developed to reduce the computational load. One of the latest in the literature is the *Split-Radix* FFT. The Split-Radix FFT [16] has the lowest number of multiplies and adds of any known algorithm. It achieves this efficiency by combining certain radix-2 and radix-4 butterflies, but, as a result, the classical concept of FFT stages is lost. The new structure uses a rather complicated indexing scheme, which is the price paid for the reduced multiplies/adds. Since, on the TMS320C30, multiplies/adds are not more expensive computationally than any other operation, the indexing scheme wipes out the gains of the reduced arithmetic. Actually, an implementation of the split-radix FFT showed it to be slower than the radix-2 FFT, one of the main reasons being that the block-repeat structure could no longer be used effectively.

Very often, there is a question on what the different benchmark numbers mean. A useful comparison of execution times for different algorithms on different machines has been made [17]. Table 4 presents a small segment of the resulting information that is relevant to the present discussion: the timing in seconds for the radix-8, mix-radix, and split-radix algorithms that were implemented on various machines. Different operating systems and compilers have been used, as shown. The execution times of Table 4 should be compared with the 0.0010055 s that it takes to implement a 1024-point, radix-2, real FFT on a TMS320C30. As can be seen, the TMS320C30 compares favorably to all the other machines investigated.

**Table 4. Execution Times in Seconds for a 1024-Point Real FFT. The Numbers Should Be Compared with 0.001055 s of a 1024-Point Real FFT on the TMS320C30**

Machine	Radix-8	Mix-radix	Split-radix
VAX 750 UNIX BSD4.2 f77	0.3634	0.3902	0.3021
VAX 750 UNIX BSD4.2 f77 -O	0.2376	0.2948	0.2089
VAX 750 UNIX BSD4.3 f77	0.2545	0.2600	0.2371
VAX 750 UNIX BSD4.3 f77 -O	0.1825	0.2127	0.1672
VAX 785 ULTRIX f77	0.1046	0.1107	0.1101
VAX 785 ULTRIX f77 -O	0.0796	0.0943	0.0811
VAX 785 VMS FOR/NOOPTM	0.0767	0.0871	0.0975
VAX 785 VMS FOR/OPTM	0.0539	0.0641	0.0633
VAX 8600 VMS FOR/OPTM	0.0217	0.0243	0.0235
MICROVAX VMS FOR/NOOPTM	0.1671	0.1846	0.1864
MICROVAX VMS FOR/OPTM	0.1299	0.1527	0.1419
DEC-10 TOPS-10 FOR/NOOPTM	0.0940	0.1184	0.0991
DEC-10 TOPS-10 FOR/OPTM	0.0885	0.1110	0.0845
CDC 855 FTN5,OPT=0	0.0277	0.0319	0.0338
CDC 855 FTN5,OPT=1	0.0277	0.0316	0.0337
CDC 855 FTN5,OPT=2	0.0182	0.0171	0.0151
CDC 855 FTN5,OPT=3	0.0180	0.0173	0.0150
SUN 3/50 UNIX BSD4.2 f77 -O -f68881	0.2518	0.3365	0.2103
SUN 3/50 UNIX BSD4.2 f77 -f68881	0.2806	0.3897	0.2802
SUN 3/50 UNIX BSD4.2 f77 -O	0.7586	1.047	0.6955
SUN 3/50 UNIX BSD4.2 f77	0.7476	1.029	0.7033
SUN 3/160 UNIX BSD4.2 f77	0.6037	0.6895	0.5660
SUN 3/160 UNIX BSD4.2 f77 -pfa	0.0983	0.1060	0.0946
SUN 3/260 UNIX BSD4.3 f77	0.3689	0.4126	0.3390
SUN 3/260 UNIX BSD4.3 f77 -O	0.3530	0.4142	0.3297
Pyramid 90X UNIX BSD4.2 f77 -O	0.2053	0.2244	0.1416
Pyramid 90X UNIX BSD4.2 f77	0.2206	0.2457	0.1326
HP-1000 21MX-E FTN7X	0.9400	1.248	0.9478
Apple MAC Microsoft FOR	2.6670	3.1600	2.8260
AST PC Microsoft FOR	1.5040	2.0800	1.4630

## The TMS320C30 C Compiler

The C compiler for the TMS320C30 permits easy porting of high-level language programs to the DSP device. If the CPU loading of a particular application is not very high, the C compiler can create programs that run on the TMS320C30 in real time. If, however, the result is non-realtime, it may be necessary to use assembly language for more efficient coding.

In most cases, only a portion of the code needs to be written in assembly language. Typically, there are a few code segments where the device spends most of the time and which, when optimized in assembly language, yield the necessary performance improvement. By following the conventions outlined in the run-time environment of the C compiler [15], you can write these time-critical routines in assembly language and call them in a C program. This is also true for the FFT routines. In appendices A, B, and C, the radix-2, radix-4, and real FFT routines mentioned earlier are also put in a C-callable form by adding the necessary interface at the beginning and the end of the code. The tables with the sines and cosines are again assumed to be supplied during link time.

### Issues in FFT Implementation

There are many ways of actually implementing the FFT code (and the other transformations), taking into consideration the different possibilities of program locations, the data locations, the ways of input and output, etc. Since it is impractical to cover every possible case, this report has concentrated on a configuration in which the use of external memory is minimized. With the source code and additional explanations provided, you should be able to customize the FFT implementation for a particular application.

### Use of External Memory

In these implementations, only on-chip memory was used, and that's why the maximum transform size considered was 1024 points long (2048 for a real transform). Often, though, applications call for use of external memory for program or data or both. When external memory is used, the structure of the code does not change at all; it is only the timing that may be affected.

Fast external memory can be selected so that no wait states are necessary. But even when there are no wait states, accessing external memory may impose some limitations. For instance, you can make only one external memory access in a full cycle, but you can make two accesses of internal memory in each cycle. Also, because of multiplexing of the busses, pipeline conflicts may arise if both program and data are placed on the same external port. Resolution of such conflicts causes extra cycles for the execution. The section on pipelining in the *TMS320C30 User's Guide* explains in detail what kind of potential conflicts may occur.

To minimize or avoid such conflicts, there are some simple steps that the designer can take. The TMS320C30 has three separate memory areas (one on-chip, one accessed by the primary bus, and one accessed by the expansion bus) that can be combined. For instance, the program can be placed on the expansion port and the data on the primary port. Or the data can first be brought into internal memory and then operated upon. Alternatively, the program may be relocated to internal memory. A related approach is to use the cache. All the transforms are implemented as loops that are executed many times. If you activate the on-chip cache after the first access of the code, the instructions execute from the cache instead of the external memory.

If there are additional conflicts, they can typically be resolved by some rearrangement of the code. For instance, consecutively writing to external memory takes two cycles per write. If, however, a write is followed by some internal operation, then the second cycle of the write is transparent, and the actual cost is one cycle.

## Bit Reversal

The TMS320C30 has a special form of the indirect addressing mode for the bit-reversing operation that is required at the beginning or the end of an FFT. Through this addressing mode, the scrambled data are accessed in their proper order. This addressing mode works as follows:

Let  $AR_n$  ( $n=0..7$ ) be the auxiliary register pointing to the array with scrambled data. The index register  $IR_0$  contains a number equal to one-half the size of the FFT. Then, after every access of the data,  $AR_n$  is incremented by  $IR_0$  using the construct

$$*AR_n + + [IR_0]B$$

This causes the contents of  $AR_n$  to be incremented by the contents of  $IR_0$ , but if there is a carry in this incrementing, the carry propagates to the right instead of to the left. The result is the generation of the addresses in a bit-reversed order. The bit-reversed addressing mode works correctly if the array with the data is aligned in memory so that the first memory address is a multiple of the FFT size. This can be achieved if the first memory address has zeros for the last  $M$  bits, where  $M = \log_2 N$ , with  $N$  being the FFT size. For example, in the case of a 1024-point FFT, the last 10 bits of the memory address of the first datum should be zeros.

In the implementation of the complex FFT, the output is complex even when the input is real. So, there is a need to consider both the real and the imaginary parts of the data array. The above description of the bit-reversed addressing mode assumed that the real and the imaginary parts are stored as separate arrays in the memory. In this case, each of the arrays (real or imaginary parts) can be accessed as described. However, in most cases (including this report), the real and imaginary points alternate in the same array.

In this arrangement, the following simple modification achieves the same goal: set IRO equal to  $N$  instead of  $N/2$ , and access the  $N$  points of the transform. At every access, the auxiliary register is pointing to the real part of the FFT. The imaginary part is located in the next higher location, and it can be easily accessed.

With the bit-reversed addressing mode, the unscrambling of the data can take place when the FFT result is accessed for further processing or for I/O. It is possible, though, that certain applications demand the reordering of the data in the same array. Such a rearrangement can be done very simply for a complex FFT with the following code.

; DO THE BIT-REVERSING EXPLICITLY

```

LDI  @FFTSIZ,RC      ; RC = FFT SIZE
SUBI 1,RC            ; RC SHOULD BE ONE LESS THAN DESIRED #
LDI  @FFTSIZ,IRO     ; IRO = FFT SIZE
LDI  @INPUT,ARO      ;
LDI  @INPUT,AR1      ;
*
RPTB BITRV
CMPI AR1,ARO         ; EXCHANGE LOCATIONS ONLY
BGE  CONT            ; IF AROAR1
LDF  *ARO,RO         ;
||   LDF  *AR1,R1     ; EXCHANGE REAL PARTS
STF  RO,*AR1        ;
||   STF  R1,*ARO     ;
LDF  *+ARO,RO        ;
||   LDF  *+AR1,R1    ; EXCHANGE IMAGINARY PARTS
STF  RO,*+AR1       ;
||   STF  R1,*+ARO    ;
CONT NOP *ARO++[2]
BITRV NOP *AR1++[IRO]B

```

Note that AR1 is pointing to the bit-reversed version of the address contained in ARO. For real-valued FFT, or for FFTs that store the real and the imaginary parts in separate arrays, the real-FFT routine in Appendix C contains a modified example of the above code.

## Use of DMA

If the signal to be transformed arrives as a continuous stream of data, the DMA could be used to collect the new data while the data already collected are processed. In this case, the data source address of the DMA points to the memory location corresponding to a serial port, or to another port associated with an external device. The destination is a memory space designated for storage.

There are two ways to use such buffers. One possibility is to designate one buffer as the temporary storage and the other buffer as the working area. When the storage buffer receives the necessary amount of data, the data is transferred to the working area, and the DMA starts refilling the storage buffer. Alternatively, the two buffers are considered equivalent: when the processor finishes processing and outputting the data from one and the DMA has filled the other, the two buffers switch functions; i.e., the DMA starts filling the first buffer while the CPU is processing the data in the buffer just filled.

## **Test Vector**

For testing purposes, a vector with 64 (quasi-random) data points and the corresponding FFT values is given in Appendix F. In this way, if any of the routines is implemented, the test vectors can be used to verify the correct functionality of the routines. Together with the test vectors, Appendix C gives a sine/cosine table for a 64-point transform, and the linking file for such a transform.

## **Summary**

This report examined implementations of fast transforms on the Texas Instruments TMS320C3x floating-point devices. The transforms considered were several forms of the FFT, the Discrete Hartley Transform, and the Discrete Cosine Transform. Because of the powerful architecture of the device, the implementation was done easily and efficiently. It was shown that a TMS320C30 executes the FFTs several times faster than large computers such as VAX and SUN workstations. With the availability of the C compiler, these routines can be put in C-callable form and be used to compute the corresponding transforms efficiently.

## Appendices

Appendices A to F contain the TMS320C30 assembly language programs for the different algorithms considered. The contents of the appendices are as follows:

### Appendix A: Radix-2 Complex FFT.

composed of

- A1: Generic Program to Do a Looped-Code Radix-2 FFT Computation on the TMS320C30.
- A2: `fft_2` - Radix-2 Complex FFT to Be Called as a C Function.
- A3: Complex, Radix-2 DIT FFT - R2DIT.ASM.
- A4: Complex, Radix-2 DIT FFT - R2DITB.ASM.
- A5: TWID1KBR.ASM - Table with Twiddle Factors for a FFT up to a Length of 1024 Complex Points.

### Appendix B: Radix-4 Complex FFT.

composed of

- B1: Generic Program to Do a Looped-Code Radix-4 FFT on the TMS320C30.
- B2: `fft_4` - Radix-4 Complex FFT to Be Called as a C Function.

### Appendix C: Radix-2 Real FFT.

composed of

- C1: Generic Program to Do a Radix-2 Real FFT Computation on the TMS320C30.
- C2: `fft_rl` - Radix-2 Real FFT to Be Called as a C Function.
- C3: Generic Program to Do a Radix-2 Real Inverse FFT Computation on the TMS320C30.

### Appendix D: Discrete Hartley Transform.

composed of

- D1: Generic Program to Do a Radix-2 Hartley Transform on the TMS320C30.

### Appendix E: Discrete Cosine Transform.

composed of

- E1: A Fast Cosine Transform.
- E2: A Fast Cosine Transform (Inverse Transform).
- E3: FCT Cosine Tables File.
- E4: Data File.

Appendix F: Test Vectors, 64-Point Sine Table, Link Command File.  
composed of

- F1: Example of a 64-Point Vector to Test the FFT Routines.
- F2: File to Be Linked with the Source Code for a 64-Point, Radix-4 FFT.
- F3: Link Command File.

The first three appendices contain the code for the radix-2, complex radix-4, and real radix-2 FFT transformations. These routines are given in both the regular form and in a C-callable form. Furthermore, the contents of a file with the twiddle factors are given, as well as an example of a link command file for a 64-point FFT. Note that the source code of these routines can be downloaded from the TI DSP bulletin board (BBS) by calling (713) 274-2323. For questions regarding the BBS, call the TI DSP hotline at (713) 274-2320.

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Mr. Raimund Meyer and Mr. Karl Schwarz (Lehrstuhl für Nachrichtentechnik, University of Erlangen) provided the fast routines of Appendix A to do 1024-point, radix-2, DIT FFT. Mr. Paul Wilhelm of the University of Washington provided the routines for the Fast Cosine Transform (FCT) together with the related explanations and the test vector in Appendix E. Their contributions are gratefully acknowledged.



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## Appendix A. Radix-2 Complex FFT

# Appendix A1. Generic Program to Do a Looped-Code Radix-2 FFT Computation on the TMS320C30

```

* GENERIC PROGRAM TO DO A LOOPED-CODE RADIX-2 FFT COMPUTATION ON THE
* TMS320C30.
*
* THE PROGRAM IS TAKEN FROM THE BARRIS & PARKS BOOK, P. 111. THE (COMPLEX)
* DATA RESIDE IN INTERNAL MEMORY. THE COMPUTATION IS DONE IN-PLACE, BUT THE
* RESULT IS MOVED TO ANOTHER MEMORY SECTION TO DEMONSTRATE THE BIT-REVERSED
* ADDRESSING. THE TWIDDLE FACTORS ARE SUPPLIED IN A TABLE PUT IN A .DATA
* SECTION. THIS DATA IS INCLUDED IN A SEPARATE FILE TO PRESERVE THE GENERIC
* NATURE OF THE PROGRAM. FOR THE SAME PURPOSE, THE SIZE OF THE FFT N AND
* LOG2(N) ARE DEFINED IN A .GLOBAL DIRECTIVE AND SPECIFIED DURING LINKING.
*
* AUTHOR: PAWS E. PAPPACHALIS
* TEXAS INSTRUMENTS
*
* JULY 16, 1987
*
* .GLOBAL FFT
* .GLOBAL N
* .GLOBAL LOG2(N)
* .GLOBAL SINE
*
* .USECT "IN",1024
* .BSS OUTP,1024
*
* .TEXT
*
* INITIALIZE
*
* .WORD FFT
*
* .SPACE 100
*
* FFTSIZ .WORD N
* LOGFT .WORD N
* SINTAB .WORD SINE
* INPUT .WORD INP
* OUTPUT .WORD OUTP
*
* FFT: LIP
* LDI @FFTSIZ,IR1
* LSH -2,IR1
* LDI @FFTSIZ,IR0
* LSH 1,IR0
* LDI @FFTSIZ,R7
* R7=MC
* LDI 1,AMS
*
* LDI 1,AMS
*
* OUTER LOOP
*
* LOOP: NOP
* LDI @INPUT,ARO
*
* ; ENTRY POINT FOR EXECUTION
* ; FFT SIZE
* ; LOG2(N)
* ; ADDRESS OF SINE TABLE
* ; MEMORY WITH INPUT DATA
* ; MEMORY WITH OUTPUT DATA
*
* ; STARTING LOCATION OF THE PROGRAM
* ; RESERVE 100 WORDS FOR VECTORS, ETC.
*
* ; COMMAND TO LOAD DATA PAGE POINTER
*
* ; TR1=N/4, POINTER FOR SIN/COS TABLE
* ; AR6 HOLDS THE CURRENT STAGE NUMBER
*
* ; IR0=2*MI (BECAUSE OF REAL/IMAG)
* ; INITIALIZE REPEAT COUNTER OF FIRST
* ; LOOP
* ; INITIALIZE IE INDEX (AR6=IE)
*
* ; CURRENT FFT STAGE
* ; ARO POINTS TO X(L)
*
* ; AR2 POINTS TO X(L)
* ; RC SHOULD BE ONE LESS THAN DESIRED #
*
* ; FIRST LOOP
*
* R7,ARO,AR2
* LDI AR7,RC
* SUBI 1,RC
*
* ; AR2 POINTS TO X(L)
* ; RC SHOULD BE ONE LESS THAN DESIRED #
*
* RPTB BLK1
* ADF @R2,++AR2,RO
* SUB @R2,++AR2,++R1
* ADF @R2,++AR2,RO
* SUB @R2,++AR2,++R2
* ADF @R2,++AR2,RO
* SUB @R2,++AR2,++R3
* ST R2,++AR2
* ST R3,++AR2
*
* BLK1 ST R0,++AR0,++(LR0)
* ST R1,++AR2,++(LR0)
*
* ; IF THIS IS THE LAST STAGE, YOU ARE DONE
*
* CMPI @LOGFT,AR6
* BZD END
*
* ; MAIN INNER LOOP
*
* LDI 2,ARI
* @SINTAB,ARA
* ADR ARA,ARA
* LDI ARI,ARO
* ADFI 2,ARI
* ADFI @INPUT,ARO
* ADFI R7,ARO,AR2
* LDI AR7,RC
* SUBI 1,RC
* LDF ++AR4,R6
*
* ; INIT LOOP COUNTER FOR INNER LOOP
* ; INITIALIZE IA INDEX (AR4=IA)
* ; IA=IA+IE; AR4 POINTS TO COSINE
* ; INCREMENT INNER LOOP COUNTER
* ; (X(L),Y(L)) POINTER
* ; (X(L),Y(L)) POINTER
* ; RC SHOULD BE ONE LESS THAN DESIRED #
* ; R6=SIN
*
* BLK2 BLK2
* ADF @R2,++AR2,R2
* SUB @R2,++AR2,R1
* MOVF R2,R6,RO
* ADF @R2,++AR2,R3
* MOVF R1,++AR4(IR1),R3
* ST R3,++AR2
* ST R4,RO
* MOVF R1,R4,RO
* ADF @R2,++AR2,R3
* MOVF R2,++AR4(IR1),R3
* ADF R3,++AR2,++(LR0)
* ST R3,++AR2,++(LR0)
*
* BLK2 ST R5,++AR2,++(LR0)
* ST R4,++AR2
*
* CMPI R7,ARI
* BNE INL0P
*
* ; LOOP BACK TO THE INNER LOOP

```

```
*
*
*     LSH      1,AR7      ; INCREMENT LOOP COUNTER FOR NEXT TIME
*     LSH      1,AR5      ; IE=2*IE
*     LODI     R7,IRO      ; RI=RI
*     LSH      -1,R7       ; R2=R2/2
*     BR       LOOP
*
*
*     STORE RESULT OUT USING BIT-REVERSED ADDRESSING
*
*     LODI     @FTSIZ,RC   ; RC=M
*     SUBI     1,RC        ; RC SHOULD BE ONE LESS THAN DESIRED #
*     LODI     @FTSIZ,IRO  ; IRO=SIZE OF FFT*M
*     LODI     2,IR1
*     RIMPUT,ARO
*     LODI     @OUTPUT,ARI
*
*     RPTB
*     LDF      ++ARO(1),RO
*     LDF      ++ARO++(IRO)B,R1
*     STF      RO,++ARI(1)
*     STF      R1,++ARI++(IR1)
*
*     SELF    BR         SELF      ; BRANCH TO ITSELF AT THE END
*     .END
```

# Appendix A2. fft\_2 - Radix-2 Complex FFT to Be Called as a C Function

```

* NAME:
* fft_2 --- RADIX-2 COMPLEX FFT TO BE CALLED AS A C FUNCTION.
*
* SYNOPSIS:
* INT fft_2(N, R, DATA)
* INT N      FFT SIZE; N=2**M
* INT M      NUMBER OF STAGES = LOG2(N)
* FLOAT *DATA ARRAY WITH INPUT AND OUTPUT DATA
*
* DESCRIPTION:
* GENERIC FUNCTION TO DO A RADIX-2 FFT COMPUTATION ON THE ZOOMCO.
* THE DATA ARRAY IS 2**M-LONG, WITH REAL AND IMAGINARY VALUES ALTERNATING.
* THE PROGRAM IS BASED ON THE FORTRAN PROGRAM IN THE BURRIS AND PARKS
* BOOK, P. 111.
*
* THE COMPUTATION IS DONE IN PLACE, AND THE ORIGINAL DATA IS DESTROYED.
* BIT REVERSAL IS IMPLEMENTED AT THE END OF THE FUNCTION. IF THIS IS NOT
* NECESSARY, THIS PART CAN BE COMMENTED OUT.
*
* THE SINE/COSINE TABLE FOR THE MIDDLE FACTORS IS EXPECTED TO BE SUPPLIED
* DURING LINK TIME, AND IT SHOULD HAVE THE FOLLOWING FORMAT:
*
* .GLOBAL   _sine
* .DATA
* _sine     .FLOAT VALUE1 = sin(0.25pi/N)
*           .FLOAT VALUE = sin(1.25pi/N)
*           .FLOAT VALUE(5M/4) = sin(1.5Mpi/4-1)25pi/N)
*
* THE VALUES VALUE1, VALUE2, ETC., ARE THE SAME WAVE VALUES, FOR AN
* N-POINT FFT, THERE ARE N**M/4 VALUES FOR A FULL AND A QUARTER PERIOD OF
* THE SINE WAVE. IN THIS WAY, A FULL SINE AND COSINE PERIOD ARE AVAILABLE
* (SUPERIMPOSED).
*
* STACK STRUCTURE UPON THE CALL:
*
*   +-----+
*   |FP(4)  | DATA
*   |FP(3)  | N
*   |FP(2)  | M
*   |FP(1)  | RETURN ADDR
*   |FP(0)  | OLD FP
*   +-----+
*
* REGISTERS USED: R0, R1, R2, R3, R4, R5, R6, R7, AR0, AR1, AR2, AR4, AR5
*                AR6, AR7, IFO, IRI, RS, RE, RC
*
* AUTHOR: PANOS E. PAPANICHAELIS
*        TEXAS INSTRUMENTS
*        OCTOBER 13, 1987
*
* *****

```

```

*
* .set      AR3
*
* .GLOBAL  _fft_2
* .GLOBAL  _sine
*
* .BSS    FETSIZ,1
* .BSS    LOGFFT,1
* .BSS    INPUT,1
*
* .TEXT
*
* .word   _sine
*
* SINTAB
*
* INITIALIZE C FUNCTION
*
* _fft_2:  PUSH    FP
*         LODI   SP,FP
*         PUSH  R4
*         PUSH  R5
*         PUSH  R6
*         PUSH  R7
*         PUSH  AR4
*         PUSH  AR5
*         PUSH  AR6
*         PUSH  AR7
*
*         LODI   R0,FP(2),R0
*         STI   R0,FFTSIZ
*         LODI   R0,FP(3),R0
*         STI   R0,LOGFFT
*         LODI   R0,FP(4),R0
*         STI   R0,RINPUT
*
*         INITIALIZE FFT ROUTINE
*
*         LODI   IFO,FFTSIZ,IRI
*         LSH   -2,IRI
*         LODI   IFO,AR6
*         IFO,FFTSIZ,IRO
*         LSH   1,IRO
*         LODI   IFO,FFTSIZ,R7
*         LODI   IFO,1,AR7
*
*         LODI   IFO,1,AR5
*
*         OUTER LOOP
*
*         NOP
*         LODI   R2,INPUT,AR0
*         ADDI  R7,AR0,AR2
*         LODI  R7,AR7,RC
*         SUBI  1,RC
*
*         LOOP:
*
*         *****

```

```

* ENTRY POINT FOR EXECUTION
* : ADDRESS OF SINE TABLE
*
* :
*
* : SAVE DEDICATED REGISTERS
*
* :
*
* : NONE ARGUMENTS TO LOCATIONS MATCHING
* : THE MARKS IN THE PROGRAM
*
* :
*
* : IRI=M/4, POINTER FOR SIN/COS TABLE
* : AR6 HOLDS THE CURRENT STAGE NUMBER
*
* : IRO=2*M1 (BECAUSE OF REAL/IMAG)
* : R7=RC
* : INITIALIZE REPEAT COUNTER OF FIRST
* : LOOP
* : INITIALIZE IE INDEX (AR5+IE)
*
* :
*
* : CURRENT FFT STAGE
* : AR0 POINTS TO X(L)
* : AR2 POINTS TO X(L)
*
* : RC SHOULD BE ONE LESS THAN DESIRED #

```

```

* FIRST LOOP
RPTB BLK1
  ADDF *AR0, *AR2, R0
  SUBF *R0-1(1)+X(L)
  SUBF *R1-1(1)-X(L)
  ADDF *AR2+, *AR0+, R1
  SUBF *R2-1(1)+Y(L)
  ADDF *AR2, *AR0, R2
  SUBF *R3-1(1)-Y(L)
  R2, *AR0-
  R3, *AR2-
  Y(L)-R3
  R0, *AR0+(1R0)
  R1, *AR2+(1R0)
  X(L)-R1 AND *AR0, 2 = *AR0, 2 + 2*H1
  X(L)-R1 AND *AR0, 2 = *AR0, 2 + 2*H1
  Y(L)-R1 AND *AR0, 2 = *AR0, 2 + 2*H1
  Y(L)-R1 AND *AR0, 2 = *AR0, 2 + 2*H1

* IF THIS IS THE LAST STAGE, YOU ARE DONE
  CPTI *LDRFTT, AR6
  BTD END

* MAIN INNER LOOP
  LDI 2, AR1
  LDI @SINTAB, AR4
  ADDI AR5, AR4
  LDI AR1, AR0
  ADDI 2, AR1
  ADDI @INPUT, AR0
  ADDI R7, AR0, AR2
  LDI AR7, RC
  SUBI 1, RC
  LDF *AR4, R6

  INI: LOOP:
  LDI @SINTAB, AR4
  ADDI AR5, AR4
  LDI AR1, AR0
  ADDI 2, AR1
  ADDI @INPUT, AR0
  ADDI R7, AR0, AR2
  LDI AR7, RC
  SUBI 1, RC
  LDF *AR4, R6

  INIT LOOP COUNTER FOR INNER LOOP
  INITIALIZE IA INDEX (AR4=IA)
  IA=IA+1E; AR4 POINTS TO COSINE

  INCREMENT INNER LOOP COUNTER
  X(L), Y(L) POINTER
  X(L), Y(L) POINTER

  RC SHOULD BE ONE LESS THAN DESIRED *
  R6-SIN

  POP AR7
  POP AR6
  POP AR5
  POP AR4
  POP R7
  POP R6
  POP R5
  POP R4
  POP R3
  POP R2
  POP R1
  RETS

* RESTORE THE REGISTER VALUES AND RETURN
  RPTB BITRV
  CPTI AR0, AR1
  BBE CONT
  LDF *AR0, R0
  LDF *AR1, R1
  LDF *AR1, R1
  LDF *AR1, R1
  LDF *AR0(1), R0
  LDF *AR0(1), R0
  LDF *AR0(1), R1
  LDF *AR1(1), R1
  LDF *AR1(1), R1
  LDF *AR1(1), R1
  LDF *AR0(2)
  NOP
  BITRV *AR1++(1R0)B

* DO THE BIT-REVERSING OF THE OUTPUT
  LDI @FFTSIZ, RC
  SUBI 1, RC
  LDI @FFTSIZ, IRO
  LDI @INPUT, AR0
  LDI @INPUT, AR1

  RPTB BITRV
  CPTI AR0, AR1
  BBE CONT
  LDF *AR0, R0
  LDF *AR1, R1
  LDF *AR1, R1
  LDF *AR1, R1
  LDF *AR0(1), R0
  LDF *AR0(1), R0
  LDF *AR0(1), R1
  LDF *AR1(1), R1
  LDF *AR1(1), R1
  LDF *AR0(2)
  NOP
  BITRV *AR1++(1R0)B

* RESTORE THE REGISTER VALUES AND RETURN
  POP AR7
  POP AR6
  POP AR5
  POP AR4
  POP R7
  POP R6
  POP R5
  POP R4
  POP R3
  POP R2
  POP R1
  RETS

```

# Appendix A3. Complex, Radix-2 DIT FFT - R2DIT.ASM

COMPLEX, RADIX-2 DIT FFT : R2DIT.ASM

GENERIC PROGRAM FOR A FAST LOOPED-CODE RADIX-2 DIT FFT COMPUTATION  
ON THE TRIS02C30

WRITTEN BY: RAHMOUD MEYER, KARL SCHWARZ 19-07-89  
LEHRSTUHL FUER WICHTDRICHTERTECHNIK  
UNIVERSITÄT ERLANGEN-NÜRNBERG  
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THE (COMPLEX) DATA RESIDE IN INTERNAL MEMORY, THE COMPUTATION IS DONE  
IN-PLACE, BUT THE RESULT IS MOVED TO ANOTHER MEMORY SECTION TO  
DEMONSTRATE THE BIT-REVERSED ADDRESSING.

FOR THIS PROGRAM THE MINIMUM FFTLENGTH IS 32 POINTS BECAUSE OF THE  
SEPARATE STAGES.

FIRST TWO PHASES ARE REALIZED AS A FOUR BUTTERFLY LOOP SINCE THE  
MULTIPLIES ARE TRIVIAL, THE MULTIPLIER IS ONLY USED FOR A LOAD IN  
PARALLEL WITH AN ADDIF OR SUBF.

EXAMPLE FOR A 1024-POINT FFT (EXCLUDING BIT REVERSAL):

MEMORY SIZE: 229 WORDS  
PROGRAM DATA (TWIDDLE FACTORS) = 512 WORDS

CYCLES PER BUTTERFLY:

STAGES 1 AND 2 = 4  
STAGES 3 TO 8 = 8  
STAGE 9 = 8.25  
STAGE 10 = 8.5

AVERAGE CYCLES/BUTTERFLY = 7.275

TOTAL BUTTERFLY/CYCLES = 37248

INITIALIZATION OVERHEAD = 2181 = 5.5% OF TOTAL TIME

TOTAL NUMBER OF INSTRUCTION CYCLES = 39429

TOTAL TIME FOR A 1024-POINT FFT = 2.36 ms (EXCLUDING BIT  
REVERSAL)

THIS PROGRAM INCLUDES FOLLOWING FILES:

THE FILE 'TWIDDKR.ASM' CONSISTS OF TWIDDLE FACTORS

THE TWIDDLE FACTORS ARE STORED IN BIT-REVERSED ORDER AND WITH A TABLE  
LENGTH OF N/2 (N = FFTLENGTH).

EXAMPLE: SHOWN FOR N=32,  $WN(n) = \cos(2\pi P1n/N) - j\sin(2\pi P1n/N)$

ADDRESS COEFFICIENT

0  $RWN(0) = \cos(2\pi P140/32) = 1$   
1  $-jUN(0) = \sin(2\pi P140/32) = 0$   
2  $RWN(4) = \cos(2\pi P144/32) = 0.707$   
3  $-jUN(4) = \sin(2\pi P144/32) = 0.707$

12  $RWN(3) = \cos(2\pi P143/32) = 0.831$

13  $-jUN(3) = \sin(2\pi P143/32) = 0.556$

14  $RWN(7) = \cos(2\pi P147/32) = 0.195$

15  $-jUN(7) = \sin(2\pi P147/32) = 0.981$

WHEN GENERATED FOR A FFT LENGTH OF 1024, THE TABLE IS FOR ALL  
AVAILABLE FFT OF LESS OR EQUAL LENGTH.

THE MISSING TWIDDLE FACTORS  $UN(1), UN(2), \dots$  ARE GENERATED BY USING  
THE SYMMETRY  $UN(N/4+n) = -jUN(n)$ . THIS CAN BE EASILY REALIZED BY  
CHANGING REAL- AND IMAGINARY PART OF THE TWIDDLE FACTORS AND BY  
NEGATING THE NEW REAL PART.

TO CHANGE THE FFT LENGTH, ONLY THE PARAMETERS IN THE HEADER OF  
TWIDDKR.ASM AND THE INPUT AND OUTPUT VECTOR LENGTHS NEED TO BE  
ALTERED.

AR + j AI ----- AR' + j AI'

+

+

+

+

+

+

+

+

+

+

+

+

+

+

+

+

+

+

+

+

+

+

+

+

+

BR + j BI ----- ( COS - j SIN ) ----- BR' + j BI'

TR = BR \* COS + BI \* SIN

TI = BR \* SIN - BI \* COS

AR = AR + TR

AI = AI - TI

BR = BR - TR

BI = BI + TI



• FIRST 2 STAGES AS RADI1-4 BUTTERFLY

• FILL PIPELINE

```

: R4 = AR + CR
: R5 = AR - CR
: R6 = DR + BR
: R7 = DR - BR
: AR' = R0 + R4 + R6
: R1 = D1, BR' = R3 = R4 - R6
: R0 = B1 + D1, AR' = R0
: R1 = B1 - D1, BR' = R3
: CR' = R2 = R5 + R1
: R1 = C1, DR' = R3 = R5 - R1
: R2 = A1 + C1, CR' = R2
: R6 = A1 - C1, DR' = R3
: A1' = R4 = R2 + R0

```

• RADI1-4 BUTTERFLY LOOP

```

BLK1
RPTB
PPYF #R2-, #R7, R0
SUBF R0, R2, R2
PPYF #R1+, #R7, R1
ADDF R7, R6, R3
ADDF R0, #R0, R4
SIF R4, #R4+
SUBF R0, #R0+
SIF R2, #R2+
SUBF R7, R6, R7
ADDF R1, #R3, R6
SIF R7, #R6+
SUBF R1, #R3+, R7
SIF R3, #R2+
ADDF R6, R4, R0
PPYF #R3+, #R7, R1
SUBF R6, R4, R3
ADDF R1, #R1, R0
SIF R0, #R4+
SUBF R1, #R1+, R1
SIF R3, #R3+
ADDF #+R2, #R7, R1
SUBF R1, R3, R3
ADDF R2, #R2+
SIF R2, #R2+
SUBF R1, #R0+, R6
SIF R3, #R6+
ADDF R0, R2, R4

```

• INPUT VECTOR LENGTH = 2N (DEPENDS)

```

: ON N)
: OUTPUT VECTOR LENGTH = 2N (DEPENDS)
: ON N)

```

• LOAD PAGE POINTER

```

: IRO = N/2 = OFFSET BETWEEN INPUTS
: #R7 POINTS TO THIDDLE FACTOR 1
: #R0 POINTS TO AR
: #R1 POINTS TO BR
: #R2 POINTS TO CR
: #R3 POINTS TO DR
: #R4 POINTS TO AR'
: #R5 POINTS TO BR'
: #R6 POINTS TO DR'
: #R7 = FIRST THIDDLE FACTOR = 1
: ADDRESS OFFSET
: IRO = N/4 = NUMBER OF R4-BUTTERFLIES

```

• FFT:

```

LUP
LDI #62, IRO
LDI #SINTAB, AR7
LDI #R0, AR0
ADDI IRO, AR0, AR1
ADDI IRO, AR1, AR2
ADDI IRO, AR2, AR3
LDI AR0, AR4
LDI AR1, AR5
LDI AR3, AR6
LDI 2, IR1
LSH -1, IRO
LDI IRO, RC
SUBI 2, RC

```

•

```

BLK1  ADDF   R0,R2,R4
*
*   CLEAR PIPELINE
*
*   SUBF   R0,R2,R2
*   ADDF   R7,R6,R3
*   STF    R4,R4R4
*   STF    R2,R4R5
*   SUBF   R7,R6,R7
*   STF    R7,R4R6
*   STF    R3,R4-4R2

*   THIRD TO LAST OF STAGE 2
*

LDF    #F02,IR1
LDF    IRO,AR5
LDF    1,AR5
LDF    1,AR6

STAGE  LDI    #SINTAB,AR7
        LDI    0,ARH
        LDI    #INPUT,AR0
        LDI    AR0,AR2
        ADDI   IRO,AR0,AR3
        LDI    AR3,AR1
        LSH   1,AR6
        LSH   -2,AR5
        LSH   1,AR5
        LSH   -1,IRO

*
        LSH   -1,IR1
        ADDI   1,IR1

*
        LDF   +AR1++,R6
        LDF   +AR7,R7

*
*   GRUPPE
*
*   FILL PIPELINE
*
*
*   LDF    +++AR7,R6
*   PVVF   +AR1-,-R6,RI
*   ADDF   +++AR4,R0,R3
*   PVVF   +AR1,R7,R0
*   ADDF   R0,R1,R3
*   PVVF   +AR1++,+AR7--,R0
*   PVVF   +AR1++,R7,R1
*   SUBF   R3,+AR0,R2
*   ADDF   +AR0++,R3,R5

*

```

```

: AI' = R4 = R2 + R0
: BI' = R2 = R2 - R0
: CI' = R3 = R6 + R7
: AI' = R4, BI' = R2
: DI' = R7 = R6 - R7
: DJ' = R7, CI' = R3

: POINTER TO TIDDLE FACTOR
: GROUP COUNTER
: UPPER REAL BUTTERFLY INPUT
: UPPER REAL BUTTERFLY OUTPUT
: LOWER REAL BUTTERFLY INPUT
: LOWER REAL BUTTERFLY OUTPUT
: DOUBLE GROUP COUNT
: HALF BUTTERFLY COUNT
: CLEAR LSB
: HALF STEP FROM UPPER TO LOWER REAL
: PHRT
: STEP FROM OLD IMAGINARY TO NEW REAL
: VALUE
: DUMMY LOAD, ONLY FOR ADDRESS UPDATE
: R7 = COS

: AR0 = UPPER REAL BUTTERFLY INPUT
: AR1 = LOWER REAL BUTTERFLY INPUT
: AR2 = UPPER REAL BUTTERFLY OUTPUT
: AR3 = LOWER REAL BUTTERFLY OUTPUT
: THE IMAGINARY PART HAS TO FOLLOW
: R6 = SIN
: R1 = BI * SIN
: DUMMY ADDF FOR COUNTER UPDATE
: R0 = BR * COS
: R3 = TR * R0 + R1, R0 = BR * SIN
: R1 = BI * COS, R2 = AR - TR
: R5 = AR + TR, BR' = R2

```

```

*
*   FIRST BUTTERFLY-TYPE:
*
*   TR = BR * COS + BI * SIN
*   TI = BR * SIN - BI * COS
*   AR' = AR + TR
*   AI = AI - TI
*   BR' = AR - TR
*   BI' = AI + TI
*
*
RPTB   BFLV1
PVVF   +AR1,R6,R5
STF    R5,AR2Z++
SUBF   R1,R0,R2
PVVF   +AR1,R7,R0
SUBF   R2,+AR0++,R4
STF    R3,AR3++
ADDF   R0,R5,R3
PVVF   +AR1++,R6,R0
SUBF   R3,+AR0,R2
PVVF   +AR1++,R7,R1
STF    R4,AR2Z++
ADDF   +AR0++,R3,R5
STF    R2,AR3++

*   SWITCH OVER TO NEXT GROUP
SUBF   R1,R0,R2
ADDF   R2,+AR0,R3
STF    R5,AR2Z++
SUBF   R2,+AR0++(IR1),R4
STF    R2,+AR3++(IR1)
NDP    +AR1++(IR1)
PVVF   +AR1-,-R7,R1
STF    R4,+AR2++(IR1)
PVVF   +AR1,R6,R0
PVVF   +AR1++,+AR7++,R0
SUBF   R0,R1,R3
PVVF   +AR1++,R6,R1
SUBF   R3,+AR0,R2
ADDF   +AR0++,R3,R5
STF    R2,AR3++
LDF    IRO,AR5

*
: R5 = BI * SIN, (AR' = R5)
: (R2 = TI = R0 - R1)
: R0 = BR * COS, (R3 = AI + TI)
: (R4 = AI - TI, BI' = R3)
: R3 = TR = R0 + R5
: R0 = BR * SIN, R2 = AR - TR
: R1 = BI * COS, (AI' = R4)
: R5 = AR + TR, BR' = R2
: R2 = TI = R0 - R1
: R3 = AI + TI, AR' = R5
: R4 = AI - TI, BI' = R3
: ADDRESS UPDATE
: R1 = BI * COS, AI' = R4
: R0 = BR * SIN
: R3 = TR = R1 - R0, R0 = BR * COS
: R1 = BI * SIN, R2 = AR - TR
: R5 = AR + TR, BR' = R2

```

```

*****
* SECOND BUTTERFLY-TYPE:
*
* TR = B1 * COS - BR * SIN
* T1 = B1 * SIN + BR * COS
* AR' = AR + TR
* AI' = AI - T1
* BR' = AR - TR
* BI' = AI + T1
*
*****
*
* RPTB BFLY2
*
* MPYF ++AR1,R7,R5 ; R5 = B1 * COS , (AR' = R5)
* STP R5,HRZ2++
*
* AOUF R1,R0,R2 ; (R2 = T1 = R0 + R1)
* MPYF ++R1,R6,R0 ; R0 = BR * SIN , (R3 = AI + T1)
* AOUF R2,HR0,R3
* SUBF R2,HR0++R4 ; (R4 = AI - T1 , BI' = R3)
*
* STP R3,HRZ3++
* SUBF R0,R5,R2 ; TR = R3 = R5 - R0
* MPYF ++R1++R7,R0 ; R0 = BR * COS , R2 = AR - TR
* SUBF R3,HR0,R2
*
* MPYF ++R1++R6,R1 ; (R1 = B1 * SIN , (AI' = R4)
* STP R4,HRZ2++
* BFLY2 AOUF ++R0++R3,R5 ; R5 = AR + TR , BR' = R2
* STP R2,HRZ3++
*
*
* CLEAR PIPELINE
*
* AOUF R1,R0,R2 ; R2 = T1 = R0 + R1
* AOUF R2,HR0,R3 ; R3 = AI + T1
* STP R3,HRZ2++
*
* CMP1 HR0,HR4
* BMD GRUPE ; DO FOLLOWING 3 INSTRUCTIONS
* SUBF R2,HR0++(IR1),R4 ; R4 = AI - T1 , BI' = R3
* STP R3,HRZ3++(IR1)
*
* LDF ++HR7,R7 ; R7 = COS
* STP R4,HRZ2++(IR1) ; AI' = R4
* NOP ++R1++(IR1) ; BRANCH HERE
*
* END OF THIS BUTTERFLY GROUP
*
* CMP1 4,IR0 ; JUMP OUT AFTER LD(IN)-3 STAGE
* BNZ STUFE
*
*
* SECOND TO LAST STAGE
*
* LDI RINPUT,AR0 ; UPPER INPUT
* LDI AR0,AR2 ; UPPER OUTPUT
* AOU1 IR0,AR0,AR1 ; LOWER INPUT

```

```

LDI AR1,AR3 ; LOWER OUTPUT
LDI @SIMP2,AR7 ; POINTER TO TUPLE FACTOR
LDI 5,IR0 ; DISTANCE BETWEEN TWO GROUPS
LDI @GBR2,RC

*
* FILL PIPELINE
*
* 1. BUTTERFLY: W'0
*
* AOUF ++R0,HR1,R2 ; AR' = R2 = AR + BR
* SUBF ++R1++R0++R3 ; BR' = R3 = AR - BR
* AOUF ++R0,HR1,R0 ; AI' = R0 = AI + BI
* SUBF ++R1++R0++R1 ; BI' = R1 = AI - BI
*
*
* 2. BUTTERFLY: W'0
*
* AOUF ++R0,HR1,R6 ; AR' = R6 = AR + BR
* SUBF ++R1++R0++R7 ; BR' = R7 = AR - BR
* AOUF ++R0,HR1,R4 ; AI' = R4 = AI + BI
* SUBF ++R1++(IR0),HR0++(IR0),R5 ; BI' = R5 = AI - BI
*
* STP R2,HRZ2++ ; (AR' = R2)
* STP R3,HRZ3++ ; (BR' = R3)
* STP R0,HRZ2++ ; (AI' = R0)
* STP R1,HRZ3++ ; (BI' = R1)
* STP R6,HRZ2++ ; AR' = R6
* STP R7,HRZ3++ ; BR' = R7
* STP R4,HRZ2++(IR0) ; AI' = R4
* STP R5,HRZ3++(IR0) ; BI' = R5
*
* 3. BUTTERFLY: W'1/4
*
* AOUF ++R0++R1,R5 ; AR' = R5 = AR + BI
* SUBF ++R1,HR0,R4 ; AI' = R4 = AI - BR
* AOUF ++R1++R0--R6 ; BI' = R6 = AI + BR
* SUBF ++R1++R0++R7 ; BR' = R7 = AR - BI
*
* 4. BUTTERFLY: W'3/4
*
* AOUF ++R1,++HR0,R3 ; AR' = R3 = AR + BI
* LDF ++R7,R1 ; R1 = 0 (FOR INNER LOOP)
* LDF ++R1++R0 ; R0 = BR (FOR INNER LOOP)
* SUBF ++R1++(IR0),HR0++R2 ; BR' = R2 = AR - BI
* STP R5,HRZ2++ ; (AR' = R5)
* STP R7,HRZ3++ ; (BR' = R7)
* STP R6,HRZ3++ ; (BI' = R6)
*
* 5. TO H. BUTTERFLY:
*
* RPTB BFLZND
*
* LDF ++R7++R7 ; R7 = COS , ((AI' = R4))
* LDF R4,HRZ2++ ; R6 = SIN , (BR' = R2)
* LDF ++R7++R6
* STP R2,HRZ3++

```





# Appendix A4. Complex, Radix-2 DIT FFT – R2DITB.ASM

## APPENDIX A4

COMPLEX, RADIX-2 DIT FFT : R2DITB.ASM

GENERIC PROGRAM FOR A FAST LOOPED-CODE RADIX-2 DIT FFT COMPUTATION  
ON THE TRIS20C30

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THE (COMPLEX) DATA RESIDE IN INTERNAL MEMORY. THE COMPUTATION IS DONE  
IN-PLACE, BUT THE RESULT IS MOVED TO ANOTHER MEMORY SECTION TO  
DEMONSTRATE THE BIT-REVERSED ADDRESSING.

FOR THIS PROGRAM THE MINIMUM FFT LENGTH IS 32 POINTS BECAUSE OF  
THE SEPARATE STAGES.

FIRST TWO PASSES ARE REALIZED AS A FOUR BUTTERFLY LOOP SINCE THE  
MULTIPLIES ARE TRIVIAL. THE MULTIPLIER IS ONLY USED FOR A LOAD IN  
PARALLEL WITH AN ADD/ OR SUB.

EXAMPLE FOR A 1024-POINT FFT (WITH BIT REVERSAL) :

MEMORY SIZE :  
PROGRAM = 231 WORDS  
DATA = 512 WORDS

CYCLES PER BUTTERFLY :  
STAGES 1 AND 2 = 4  
STAGES 3 TO 8 = 8  
STAGE 9 = 8.25  
STAGE 10 = 10.5 (DUE TO EXT. MEMORY WAITS)

AVERAGE CYCLES/BUTTERFLY = 7.475  
TOTAL BUTTERFLY CYCLES = 36272  
INITIALIZATION OVERHEAD = 2185 = 5.4 % OF TOTAL TIME  
TOTAL NUMBER OF INSTRUCTION CYCLES = 40457  
TOTAL TIME FOR A 1024 POINT FFT = 2.42 ms (INCLUDING BIT  
REVERSAL)

THIS PROGRAM INCLUDES FOLLOWING FILES:-

THE FILE 'TUIDKOR.ASM' CONSISTS OF TWIDDLE FACTORS

THE TWIDDLE FACTORS ARE STORED IN BIT REVERSED ORDER AND WITH A TABLE  
LENGTH OF  $N/2$  ( $N = FFTLENGTH$ ).

EXAMPLE: SHOW FOR  $N=32$ ,  $M(16) = \cos(2\pi P16/N) - j\sin(2\pi P16/N)$

ADDRESS COEFFICIENT

0  $R(M(0)) = \cos(2\pi P16/32) = 1$

1  $-I(M(0)) = \sin(2\pi P16/32) = 0$

2  $R(M(4)) = \cos(2\pi P14/32) = 0.707$

3  $-I(M(4)) = \sin(2\pi P14/32) = 0.707$

12  $R(M(12)) = \cos(2\pi P12/32) = 0.851$

13  $-I(M(12)) = \sin(2\pi P12/32) = 0.356$

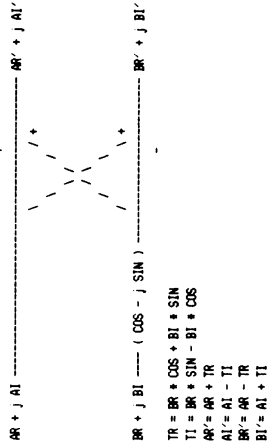
14  $R(M(17)) = \cos(2\pi P17/32) = 0.195$

15  $-I(M(17)) = \sin(2\pi P17/32) = 0.981$

WHEN GENERATED FOR A FFT LENGTH OF 1024, THE TABLE IS FOR ALL  
AVAILABLE FFT OF LESS OR EQUAL LENGTH.

THE MISSING TWIDDLE FACTORS  $(M(1), M(2), \dots)$  ARE GENERATED BY USING  
THE SYMMETRY  $M(N/4+k) = -M(N/4-k)$ . THIS CAN BE EASILY REALIZED BY  
CHANGING REAL- AND IMAGINARY PART OF THE TWIDDLE FACTORS AND BY  
NEGATING THE NEW REAL PART.

TO CHANGE THE FFT LENGTH ONLY THE PARAMETERS IN THE HEADER OF  
TUIDKOR.ASM AND THE INPUT AND OUTPUT VECTOR LENGTHS NEED TO BE  
ALTERED.



\* FIRST 2 STAGES AS RADII-4 BUTTERFLY  
 \*  
 \* FILL PIPELINE  
 \*

```

A0F2, A0R0, R4
SUBF A0C2, A0R0, R4
A0F3, A0R3, R6
SUBF A0C3, A0R3, R6
A0F4, A0R4, R8
SUBF A0C4, A0R4, R8
A0F5, A0R5, R10
SUBF A0C5, A0R5, R10
A0F6, A0R6, R12
SUBF A0C6, A0R6, R12
A0F7, A0R7, R14
SUBF A0C7, A0R7, R14
A0F8, A0R8, R16
SUBF A0C8, A0R8, R16
A0F9, A0R9, R18
SUBF A0C9, A0R9, R18
A0F10, A0R10, R20
SUBF A0C10, A0R10, R20
A0F11, A0R11, R22
SUBF A0C11, A0R11, R22
A0F12, A0R12, R24
SUBF A0C12, A0R12, R24
A0F13, A0R13, R26
SUBF A0C13, A0R13, R26
A0F14, A0R14, R28
SUBF A0C14, A0R14, R28
A0F15, A0R15, R30
SUBF A0C15, A0R15, R30
A0F16, A0R16, R32
SUBF A0C16, A0R16, R32
A0F17, A0R17, R34
SUBF A0C17, A0R17, R34
A0F18, A0R18, R36
SUBF A0C18, A0R18, R36
A0F19, A0R19, R38
SUBF A0C19, A0R19, R38
A0F20, A0R20, R40
SUBF A0C20, A0R20, R40
A0F21, A0R21, R42
SUBF A0C21, A0R21, R42
A0F22, A0R22, R44
SUBF A0C22, A0R22, R44
A0F23, A0R23, R46
SUBF A0C23, A0R23, R46
A0F24, A0R24, R48
SUBF A0C24, A0R24, R48
A0F25, A0R25, R50
SUBF A0C25, A0R25, R50
A0F26, A0R26, R52
SUBF A0C26, A0R26, R52
A0F27, A0R27, R54
SUBF A0C27, A0R27, R54
A0F28, A0R28, R56
SUBF A0C28, A0R28, R56
A0F29, A0R29, R58
SUBF A0C29, A0R29, R58
A0F30, A0R30, R60
SUBF A0C30, A0R30, R60
A0F31, A0R31, R62
SUBF A0C31, A0R31, R62
A0F32, A0R32, R64
SUBF A0C32, A0R32, R64
A0F33, A0R33, R66
SUBF A0C33, A0R33, R66
A0F34, A0R34, R68
SUBF A0C34, A0R34, R68
A0F35, A0R35, R70
SUBF A0C35, A0R35, R70
A0F36, A0R36, R72
SUBF A0C36, A0R36, R72
A0F37, A0R37, R74
SUBF A0C37, A0R37, R74
A0F38, A0R38, R76
SUBF A0C38, A0R38, R76
A0F39, A0R39, R78
SUBF A0C39, A0R39, R78
A0F40, A0R40, R80
SUBF A0C40, A0R40, R80
A0F41, A0R41, R82
SUBF A0C41, A0R41, R82
A0F42, A0R42, R84
SUBF A0C42, A0R42, R84
A0F43, A0R43, R86
SUBF A0C43, A0R43, R86
A0F44, A0R44, R88
SUBF A0C44, A0R44, R88
A0F45, A0R45, R90
SUBF A0C45, A0R45, R90
A0F46, A0R46, R92
SUBF A0C46, A0R46, R92
A0F47, A0R47, R94
SUBF A0C47, A0R47, R94
A0F48, A0R48, R96
SUBF A0C48, A0R48, R96
A0F49, A0R49, R98
SUBF A0C49, A0R49, R98
A0F50, A0R50, R100
SUBF A0C50, A0R50, R100

```

\* RADII-4 BUTTERFLY LOOP

```

BLK1
A0F2, A0R2, R2
SUBF A0C2, A0R2, R2
A0F3, A0R3, R4
SUBF A0C3, A0R3, R4
A0F4, A0R4, R6
SUBF A0C4, A0R4, R6
A0F5, A0R5, R8
SUBF A0C5, A0R5, R8
A0F6, A0R6, R10
SUBF A0C6, A0R6, R10
A0F7, A0R7, R12
SUBF A0C7, A0R7, R12
A0F8, A0R8, R14
SUBF A0C8, A0R8, R14
A0F9, A0R9, R16
SUBF A0C9, A0R9, R16
A0F10, A0R10, R18
SUBF A0C10, A0R10, R18
A0F11, A0R11, R20
SUBF A0C11, A0R11, R20
A0F12, A0R12, R22
SUBF A0C12, A0R12, R22
A0F13, A0R13, R24
SUBF A0C13, A0R13, R24
A0F14, A0R14, R26
SUBF A0C14, A0R14, R26
A0F15, A0R15, R28
SUBF A0C15, A0R15, R28
A0F16, A0R16, R30
SUBF A0C16, A0R16, R30
A0F17, A0R17, R32
SUBF A0C17, A0R17, R32
A0F18, A0R18, R34
SUBF A0C18, A0R18, R34
A0F19, A0R19, R36
SUBF A0C19, A0R19, R36
A0F20, A0R20, R38
SUBF A0C20, A0R20, R38
A0F21, A0R21, R40
SUBF A0C21, A0R21, R40
A0F22, A0R22, R42
SUBF A0C22, A0R22, R42
A0F23, A0R23, R44
SUBF A0C23, A0R23, R44
A0F24, A0R24, R46
SUBF A0C24, A0R24, R46
A0F25, A0R25, R48
SUBF A0C25, A0R25, R48
A0F26, A0R26, R50
SUBF A0C26, A0R26, R50
A0F27, A0R27, R52
SUBF A0C27, A0R27, R52
A0F28, A0R28, R54
SUBF A0C28, A0R28, R54
A0F29, A0R29, R56
SUBF A0C29, A0R29, R56
A0F30, A0R30, R58
SUBF A0C30, A0R30, R58
A0F31, A0R31, R60
SUBF A0C31, A0R31, R60
A0F32, A0R32, R62
SUBF A0C32, A0R32, R62
A0F33, A0R33, R64
SUBF A0C33, A0R33, R64
A0F34, A0R34, R66
SUBF A0C34, A0R34, R66
A0F35, A0R35, R68
SUBF A0C35, A0R35, R68
A0F36, A0R36, R70
SUBF A0C36, A0R36, R70
A0F37, A0R37, R72
SUBF A0C37, A0R37, R72
A0F38, A0R38, R74
SUBF A0C38, A0R38, R74
A0F39, A0R39, R76
SUBF A0C39, A0R39, R76
A0F40, A0R40, R78
SUBF A0C40, A0R40, R78
A0F41, A0R41, R80
SUBF A0C41, A0R41, R80
A0F42, A0R42, R82
SUBF A0C42, A0R42, R82
A0F43, A0R43, R84
SUBF A0C43, A0R43, R84
A0F44, A0R44, R86
SUBF A0C44, A0R44, R86
A0F45, A0R45, R88
SUBF A0C45, A0R45, R88
A0F46, A0R46, R90
SUBF A0C46, A0R46, R90
A0F47, A0R47, R92
SUBF A0C47, A0R47, R92
A0F48, A0R48, R94
SUBF A0C48, A0R48, R94
A0F49, A0R49, R96
SUBF A0C49, A0R49, R96
A0F50, A0R50, R98
SUBF A0C50, A0R50, R98
A0F51, A0R51, R100
SUBF A0C51, A0R51, R100

```

```

: INPUT VECTOR LENGTH = 2N (DEPENDS
: ON N)
: OUTPUT VECTOR LENGTH = 2N (DEPENDS
: ON N)

```

```

: LOAD PAGE POINTER
: IRO = N/2 = OFFSET BETWEEN INPUTS
: A0F POINTS TO TIDDLE FACTOR 1
: A0R POINTS TO AR
: A0C POINTS TO BR
: A0D POINTS TO DR
: A0E POINTS TO DR'
: A0G POINTS TO BR'
: A0H POINTS TO DR'
: ADDRESS OFFSET
: IRO = N/4 = NUMBER OF RAD-BUTTERFLIES
: RC

```

```

.global FFT
.global N
.global MHALB
.global MVIERT
.global MACHTEL
.global M
.global SINE
.bss IMP_2048
.bss OUTP_2048
.text
.word N
FFTSIZ
FGA0C2 .word MVIERT-2
FGA0C3 .word MVIERT-3
FGA0C4 .word MACHTEL-2
FGA0C5 .word MHALB-3
FGA0C6 .word M
LOCKFFT .word N
SINH1R8 .word SINE
SINH1I .word SINE-1
SINH1P2 .word SINE+2
INPUT .word IMP
OUTPUT .word IMP+2
OUTP1 .word OUTP+1
:
: A00 = AR + AI
: A01 = BR + BI
: A02 = CR + CI + CR' + CI'
: A03 = DR + DI
: A04 = AR' + AI'
: A05 = BR' + BI'
: A06 = DR' + DI'
: A07 = FIRST TIDDLE FACTOR = 1
:
: LOP
: IRO = N/2 = OFFSET BETWEEN INPUTS
: A0F POINTS TO TIDDLE FACTOR 1
: A0R POINTS TO AR
: A0C POINTS TO BR
: A0D POINTS TO DR
: A0E POINTS TO DR'
: A0G POINTS TO BR'
: A0H POINTS TO DR'
: ADDRESS OFFSET
: IRO = N/4 = NUMBER OF RAD-BUTTERFLIES
: RC
SUB1

```

```

:: STP R2, *R6++
BLK1 ANDF R0, R2, R4
* CLEAR PIPELINE
SUBF R0, R2, R2
ANDF R7, R6, R3
STF R4, *R6A
:: STP R2, *R6S
SUBF R7, R6, R7
STP R7, *R6B
STP R2, *R6C
* THIRD TO LAST-2 STAGE
LDT LDI #R2, R1
LDT LDI IRO, *R5
SUBI L, *R5
LDT LDI 1, *R6
* STAGE
LDT LDI #SIN, *R7
LDT LDI 0, *R6A
LDT LDI #INPUT, *R0
LDT LDI IRO, *R2
ANDI IRO, *R0, *R3
LDT LDI #R3, R1
LSH 1, *R6
LSH -2, *R5
LSH 1, *R5
LSH -1, IRO
*
LSH -1, R1
ANDI 1, R1
*
LDF *R1++ , R6
LDF *R7, R7
*
GRUPE
* FILL PIPELINE
LDF *R7, R6
PPVF *R1-- , R6, R1
ANDF **R6A, R0, R3
PPVF *R1, R7, R0
PPVF *R1++ , *R7-- , R0
R0, R1, R2
PPVF *R1++ , R7, R1
SUBF R3, *R0, R2
ANDF *R0++ , R2, R5
STF RZ, *R6B++
*
LDF *R1-- , R6, R1
PPVF *R1-- , R6, R1
ANDF **R6A, R0, R3
PPVF *R1, R7, R0
PPVF *R1++ , *R7-- , R0
R0, R1, R2
PPVF *R1++ , R7, R1
SUBF R3, *R0, R2
ANDF *R0++ , R2, R5
STF RZ, *R6B++

```

```

LDT AMS, R6
* FIRST BUTTERFLY-TYPE:
TR = BR * COS + B1 * SIN
TI = BR * SIN - B1 * COS
AR' = AR * TR
AI' = AI - TI
BR = AR - TR
BI' = AI * TI+
RPTB BELY1
PPVF *R1, R6, R5
STF R5, *R6C++
SUBF R1, R0, R2
PPVF *R1, R7, R0
ANDF RZ, *R0, R3
SUBF RZ, *R0++ , R4
STF R3, *R6B++
ANDF R0, R5, R3
PPVF *R1++ , R6, R0
SUBF RZ, *R0, R2
PPVF *R1++ , R7, R1
:: STP R4, *R6C++
ANDF *R0++ , R3, R5
STF RZ, *R6B3++
* SWITCH OVER TO NEXT GROUP
SUBF R1, R0, R2
ANDF RZ, *R0, R3
STF R3, *R6C++
SUBF RZ, *R0++ (R1), R4
STF RZ, *R6C++ (R1)
NOP
PPVF *R1-- , R7, R1
STF R4, *R6C++ (R1)
PPVF *R1, R6, R0
PPVF *R1++ , *R7++ , R0
SUBF R0, R1, R3
PPVF *R1++ , R6, R1
SUBF R3, *R0, R2
ANDF *R0++ , R3, R5
STF RZ, *R6B3++
LDT AMS, R6
* SECOND BUTTERFLY-TYPE:
TR = B1 * COS - BR * SIN
TI = B1 * SIN + BR * COS
AR' = AR * TR
AI' = AI - TI
BR = AR - TR

```

```

: R5 = B1 * SIN , (AR' = R5)
: (R2 = TI = R0 - R1)
: R0 = BR * COS , (R3 = AI + TI)
: (R4 = AI - TI , BI' = R3)
: R3 = TR = R0 + R5
: R0 = BR * SIN , R2 = AR - TR
: R1 = B1 * COS , (AI' = R4)
: R5 = AR + TR , BR' = R2
: R2 = TI = R0 - R1
: R3 = AI + TI , AR' = R5
: R4 = AI - TI , BI' = R3
: ADDRESS UPDATE
: R1 = B1 * COS , AI' = R4
: R0 = BR * SIN
: R3 = TR = R1 - R0 , R0 = BR * COS
: R1 = B1 * SIN , R2 = AR - TR
: R5 = AR + TR , BR' = R2

```

```

: R5 = B1 * SIN , (AR' = R5)
: (R2 = TI = R0 - R1)
: R0 = BR * COS , (R3 = AI + TI)
: (R4 = AI - TI , BI' = R3)
: R3 = TR = R0 + R5
: R0 = BR * SIN , R2 = AR - TR
: R1 = B1 * COS , (AI' = R4)
: R5 = AR + TR , BR' = R2
: R2 = TI = R0 - R1
: R3 = AI + TI , AR' = R5
: R4 = AI - TI , BI' = R3
: ADDRESS UPDATE
: R1 = B1 * COS , AI' = R4
: R0 = BR * SIN
: R3 = TR = R1 - R0 , R0 = BR * COS
: R1 = B1 * SIN , R2 = AR - TR
: R5 = AR + TR , BR' = R2

```



```

*   *   BL' = AI + TI
*
*   *   RPTB          BELY2
*
*   *   MNYF    ++R1, R7, R5        ; R5 = BI * COS, (AR' = R5)
*   *   R5, *R2++
*   *   STF     R1, R0, R2
*   *   ADF     ++R1, R6, R0
*   *   MNYF    ++R1, R6, R0
*   *   ADF     R2, ++R0, R3
*   *   STF     R3, *R2++
*   *   SUBF    R0, R5, R3
*   *   MNYF    ++R1+, R7, R0
*   *   SUBF    R3, ++R0, R2
*   *   MNYF    ++R1+, R6, R1
*   *   STF     R4, *R2++
*   *   ADF     ++R0+, R3, R5
*   *   STF     R2, *R2++
*
*   *   CLEAR PIPELINE
*   *
*   *   ADF     R1, R0, R2
*   *   ADF     R2, ++R0, R3
*   *   STF     R5, *R2++
*   *   CPE1    ARA, RRA
*   *   BNE2    GRAPE
*   *   SUBF    R2, ++R0+(IR1), R4
*   *   R4 = AI - TI, BI' = R3
*   *   STF     R3, ++R2+(IR1)
*   *   LDF     ++R7, R7
*   *   STF     R4, ++R2+(IR1)
*   *   MNP     ++R1+(IR1)
*
*   *   END OF THIS BUTTERFLY GROUP
*
*   *   CPE1    4, LPO
*   *   BNZ     STUPE
*
*   *   SECOND TO LAST STAGE
*
*   *   LDI     @INPUT, ARO
*   *   LDI     ARO, AR2
*   *   ADOI    IRO, ARO, AR1
*   *   LDI     AR1, AR3
*   *   LDI     @SUMT2, AR7
*   *   LDI     5, LPO
*   *   LDI     @COR2, RC
*
*   *   FILL PIPELINE
*
*   *   1. BUTTERFLY: W'0
*   *
*   *   ADF     ++R0, ++R1, R2
*   *   SUBF    ++R1+, ++R0+, R3
*   *   ADF     ++R0, ++R1, R0

```

```

*   *   SUBF    ++R1+, ++R0+, R1        ; BI' = R1 = AI - BI
*
*   *   2. BUTTERFLY: W'0
*   *
*   *   ADF     ++R0, ++R1, R6
*   *   SUBF    ++R1+, ++R0+, R7
*   *   ADF     ++R0, ++R1, R4
*   *   SUBF    ++R1+(IR0), ++R0+(IR0), R5        ; BI' = R5 = AI - BI
*   *   R2, *R2++
*   *   STF     R0, *R2++
*   *   STF     R1, *R2++
*   *   STF     R6, *R2++
*   *   STF     R7, *R2++
*   *   STF     R4, *R2+(IR0)
*   *   STF     R5, *R2+(IR0)
*
*   *   3. BUTTERFLY: W'1/4
*   *
*   *   ADF     ++R0+, ++R1, R5
*   *   SUBF    ++R1, ++R0, R4
*   *   ADF     ++R1+, ++R0+, R6
*   *   SUBF    ++R1+, ++R0+, R7
*
*   *   4. BUTTERFLY: W'3/4
*   *
*   *   ADF     ++R1, ++R0, R3
*   *   LDF     ++R7, R1
*   *   LDF     ++R1+(IR0), ++R0+, R2        ; BR' = R2 = AR - BI'
*   *   SUBF    R5, *R2++
*   *   STF     R7, *R2++
*   *   STF     R6, *R2++
*
*   *   5. TO M. BUTTERFLY:
*   *
*   *   RPTB    BFZND0
*
*   *   LDF     ++R7+, R7
*   *   STF     R4, *R2++
*   *   LDF     ++R7+, R6
*   *   STF     R2, *R2++
*   *   MNYF    ++R1, R6, R5
*   *   R3, *R2++
*   *   ADF     R1, R0, R2
*   *   MNYF    ++R1, R7, R0
*   *   SUBF    R2, ++R0, R3
*   *   STF     R3, *R2+(IR0), R4
*   *   ADF     R0, R5, R3
*   *   ADF     ++R1+, R6, R0
*   *   SUBF    R3, ++R0, R2
*   *   MNYF    ++R1+, R7, R1
*   *   STF     R4, *R2+(IR0)

```

```

*   *   ; R5 = BI * COS, (AI' = R4)
*   *   ; R6 = SIN, (BR' = R2)
*   *   ; R5 = BI * SIN, (AR' = R3)
*   *   ; (R2 = TI = R0 + R1)
*   *   ; R0 = BR * SIN, (R3 = AI + TI)
*   *   ; (R4 = AI - TI, BI' = R3)
*   *   ; TR = R3 = R5 - R0
*   *   ; R0 = BR * COS, R2 = AR - TR
*   *   ; R1 = BI * SIN, (AI' = R4)
*   *   ; R5 = AR + TR, BR' = R2

```

```

*   *   ; R2 = TI = R0 + R1
*   *   ; R3 = AI + TI
*   *   ; AR' = R5
*
*   *   DO FOLLOWING 3 INSTRUCTIONS
*   *   ; R4 = AI - TI, BI' = R3
*   *   ; R7 = COS
*   *   ; AI' = R4
*   *   ; BRANCH HERE
*
*   *   ; JUMP OUT AFTER LD(M)-3 STAGE
*
*   *   UPPER INPUT
*   *   ; UPPER OUTPUT
*   *   ; LOWER INPUT
*   *   ; LOWER OUTPUT
*   *   ; POINTER TO TWIDDLE FACTOR
*   *   ; DISTANCE BETWEEN TWO GROUPS
*
*   *   AR' = R2 = AR + BR
*   *   BR' = R3 = AR - BR
*   *   AI' = R0 = AI + BI

```



```

:: ANDF R2, #480, R2
PPVF #481+((IR1), R7, R1 ; R1 = BI * COS , (BI' = RA)
STF R4, #483+((IR0)B
SUBF R2, #480+*, R3 ; BR' = R3 = AR - TR , AR' = R2
:: $
STF R2, #482+((IR0)B
PPVF #481, R7, R5 ; R5 = BI * COS , (BR' = R3)
SUBF R3, #482+((IR0)B
PPVF R1, R0, R2 ; (R2 = T1 = R0 - R1)
SUBF #481, R6, R0 ; R0 = BR * SIN , (A1' = R3 = A1 - T1)
:: ANDF R2, #480, R3 ; (BI' = RA = A1 + T1 , A1' = R3)
PPVF R0, R5, R2 ; R3 = TR = R0 - R5
STF #481+*, R7, R0 ; R0 = BR * COS , AR' = R2 = AR + TR
PPVF R2, #480, R2
PPVF #481+((IR1), R6, R1 ; R1 = BI * SIN , BR' = R3 = AR - TR
SUBF R2, #480+*, R3
:: $
CLEAR PIPELINE
$
STF R2, #482+((IR0)B ; AR' = R2 , (BI' = RA)
STF R4, #483+((IR0)B ; R2 = T1 = R0 + R1
SUBF R2, #480, R3 ; A1' = R3 = A1 - T1 , BR' = R3
:: ANDF R3, #482
PPVF R2, #480, R4 ; BI' = RA = A1 + T1 , A1' = R3
STF R3, #483+((IR0)B ; BI' = RA
PPVF R4, #483
$
END OF FFT
$
END: NOP
NOP
NOP
NOP
$
SELF BR SELF
.END
$

```

# Appendix A5. TWID1KBR.ASM – Table With Twiddle Factors for a FFT up to a Length of 1024 Complex Points.

```
.float 7.11632195745216e-001
.float 7.027547445722e-001
.float 6.13288844915452e-003
.float 9.99981175282601e-001
```

```
*****
*
* APPENDIX A5
*
* TITLE: TWID1KBR.ASM
*
* TABLE WITH TWIDDLE FACTORS FOR A FFT UP TO A LENGTH OF 1024 COMPLEX
* POINTS.
*
* FILE TO BE LINKED WITH THE SOURCE CODE : K2011.ASM OR K2011B.ASM
*
* WRITTEN BY : RAIMUND NEYER AND KARL SCHWARZ
*              14.07.89
*              LEHRSTUHL FÜR INFORMATIKSYSTEMTECHNIK
*              UNIVERSITÄT DUISBURG-ESSEN
*
* LENGTH OF TWIDDLE FACTOR TABLE : 512 REAL VALUES (=1024 FFT)
*
*****
```

```
.global sine
.global n
.global nhalb
.global nqvart
.global nachtel
.global m

n .set 1024 ; FFT-LENGTH a
nhalb .set 512 ; n/2
nqvart .set 256 ; n/4
nachtel .set 128 ; n/8
m .set 10 ; NUMBER OF STAGES = 14(n)
```

```
* ANOTHER EXAMPLE OF FFT-LENGTH n = 321
* ONLY THE FIRST 16 VALUES OF THE TABLE ARE NEEDED
```

```
n .set 2
nhalb .set 16
nqvart .set 8
nachtel .set 4
m .set 5

.delta

sine
.float 1.0000000000000e+000
.float 0.0000000000000e+000
.float 7.07106781186548e-001
.float 7.07106781186548e-001
.float 9.23879532511287e-001
.float 3.82683432265990e-001
.float 3.82683432265990e-001
.float 9.23879532511287e-001
.float 9.8076280402230e-001
```

## Appendix B. Radix-4 Complex FFT



```

SUBF   *+R2, *+R0, R4
STF    R6, *+R0
SUBF   R3, R1
LUF    *+R2, R5
LUF    *+R1, R7
ADDF   *+R3, *+R1, R3
ADDF   R5, *+R0, R1
STF    R1, *+R1
ADDF   R3, R1, R6
SUBF   R5, *+R0, R2
STF    R6, *+R0+(1R0)
SUBF   R3, R1
SUBF   *+R3, *+R1, R6
SUBF   R7, *+R3, R3
SUBF   R6, R4, R5
ADDF   R6, R4
STF    R4, *+R2
STF    R3, R2, R5
SUBF   R3, R2
ADDF   R5, *+R2+(1R0)
*
*   IF THIS IS THE LAST STAGE, YOU ARE DONE
*
LDI    @STAGE, AR7
ADDI   1, AR7
CPHI   @LOFFT, AR7
BID    END
STI    AR7, @STAGE
*
*   MAIN INNER LOOP
*
LDI    1, AR7
STI    AR7, @IAI
LDI    2, AR7
STI    AR7, @LPUNT
INLDP:
LDI    2, AR6
@LPUNT, AR6
ADDI   @LPUNT, AR0
LDI    @IENDX, AR7
@INPUT, AR0
STI    AR7, @IAI
STI    AR6, @LPUNT
R0, AR0, AR1
R0, AR1, AR2
R0, AR2, AR3
@LPUNT, AR
LDI    1, AR
SUBI   @I, AR6
BID    END
*
*   CREATE COSINE INDEX, AR4
*
LDI    @IAI, AR7
@SINIB, AR4
ADDI   AR4, AR7, AR5
SUBI   1, AR5
ADDI   AR7, AR5, AR6
SUBI   1, AR6
*
*   SECOND LOOP
*
R0TB
ADDF   *+R2, *+R0, R3
ADDF   *+R3, *+R1, R5
ADDF   R5, R3, R6
SUBF   *+R2, *+R0, R4
SUBF   R5, R3
ADDF   *+R2, *+R0, R1
ADDF   *+R3, *+R1, R5
PPYF   R3, *+R5(1R1), R6
PPYF   R5, R1, R7
ADDF   R6, *+R0
SUBF   *+R2, *+R0, R2
SUBF   R5, R1
PPYF   R1, *+R5, R7
STF    R7, *+R0+(1R0)
SUBF   R7, R6
SUBF   *+R3, *+R1, R5
PPYF   R1, *+R5(1R1), R7
STF    R6, *+R1
PPYF   R3, *+R5, R6
ADDF   R7, R6
ADDF   R5, R2, R1
SUBF   R3, R2
SUBF   *+R3, *+R1, R5
ADDF   R5, R4, R3
ADDF   R3, *+R4(1R1), R6
PPYF   R1, *+R4(1R1), R6
STF    R6, *+R2
PPYF   R3, *+R4, R7
ADDF   R7, R6
PPYF   R4, *+R6(1R1), R6
PPYF   R6, *+R2+(1R0)
STF    R2, *+R6, R7
PPYF   R2, *+R6(1R1), R6
SUBF   R2, *+R6(1R1), R6
STF    R6, *+R3
STF    R4, *+R4, R7
ADDF   R7, R6
SUBI   R6, *+R3+(1R0)
STF    R6, *+R3+(1R0)
*
*   IF SHOULD BE ONE LESS THAN DESIRED
*
LDI    1, AR
SUBI   @I, AR6
BID    END

```

```

; LOOP BACK TO THE INNER LOOP
; POINT TO SIN(45)
; CREATE COSINE INDEX AREA=021

; R1=(11)*(12)
; R2=(11)-(12)
; R3=(11)*(12)
; R4=(11)-(12)
; R5=(11)*(13)
; R6=R5-R1
; R7=R5-R1
; R1=R1+R5
; R7=R3-R5
; R3=R3+R5
; Y11=R3+R5
; R11=R1+R5
; R1=(11)-(13)
; R3=(11)-(13)
; R5=(11)-R1
; Y11=R3-R5
; R1=R1+R3
; R2=R2+R3
; R3=R4-R1
; R4=R4+R1
; R1=R2-R5
; R1=R1+021
; R3=R3+R5
; R3=R3+021
; R1=R1+R2
; R4=R2-R1
; R1=R1+021
; Y11=(R3+R5)+021
; R2=R2+R4
; R2=R2+021
; Y11=(R4+R2)+021
; Y11=(R4+R2)+021
; Y11=(R4+R2)+021
; Y11=(R4+R2)+021

; LOOP BACK TO THE INNER LOOP
; INCREMENT REPEAT COUNTER FOR NEXT
; TIME
; IE=4+IE

; LOOP BACK TO THE INNER LOOP
; INCREMENT REPEAT COUNTER FOR NEXT
; TIME
; IE=4+IE

```

```

; LOOP BACK TO THE INNER LOOP
; POINT TO SIN(45)
; CREATE COSINE INDEX AREA=021
; R1=(11)*(12)
; R2=(11)-(12)
; R3=(11)*(12)
; R4=(11)-(12)
; R5=(11)*(13)
; R6=R5-R1
; R7=R5-R1
; R1=R1+R5
; R7=R3-R5
; R3=R3+R5
; Y11=R3+R5
; R11=R1+R5
; R1=(11)-(13)
; R3=(11)-(13)
; R5=(11)-R1
; Y11=R3-R5
; R1=R1+R3
; R2=R2+R3
; R3=R4-R1
; R4=R4+R1
; R1=R2-R5
; R1=R1+021
; R3=R3+R5
; R3=R3+021
; R1=R1+R2
; R4=R2-R1
; R1=R1+021
; Y11=(R3+R5)+021
; R2=R2+R4
; R2=R2+021
; Y11=(R4+R2)+021
; Y11=(R4+R2)+021
; Y11=(R4+R2)+021
; Y11=(R4+R2)+021
; LOOP BACK TO THE INNER LOOP
; INCREMENT REPEAT COUNTER FOR NEXT
; TIME
; IE=4+IE

```

```

; LOOP BACK TO THE INNER LOOP
; POINT TO SIN(45)
; CREATE COSINE INDEX AREA=021
; R1=(11)*(12)
; R2=(11)-(12)
; R3=(11)*(12)
; R4=(11)-(12)
; R5=(11)*(13)
; R6=R5-R1
; R7=R5-R1
; R1=R1+R5
; R7=R3-R5
; R3=R3+R5
; Y11=R3+R5
; R11=R1+R5
; R1=(11)-(13)
; R3=(11)-(13)
; R5=(11)-R1
; Y11=R3-R5
; R1=R1+R3
; R2=R2+R3
; R3=R4-R1
; R4=R4+R1
; R1=R2-R5
; R1=R1+021
; R3=R3+R5
; R3=R3+021
; R1=R1+R2
; R4=R2-R1
; R1=R1+021
; Y11=(R3+R5)+021
; R2=R2+R4
; R2=R2+021
; Y11=(R4+R2)+021
; Y11=(R4+R2)+021
; Y11=(R4+R2)+021
; Y11=(R4+R2)+021
; LOOP BACK TO THE INNER LOOP
; INCREMENT REPEAT COUNTER FOR NEXT
; TIME
; IE=4+IE

```

```

; LOOP BACK TO THE INNER LOOP
; POINT TO SIN(45)
; CREATE COSINE INDEX AREA=021
; R1=(11)*(12)
; R2=(11)-(12)
; R3=(11)*(12)
; R4=(11)-(12)
; R5=(11)*(13)
; R6=R5-R1
; R7=R5-R1
; R1=R1+R5
; R7=R3-R5
; R3=R3+R5
; Y11=R3+R5
; R11=R1+R5
; R1=(11)-(13)
; R3=(11)-(13)
; R5=(11)-R1
; Y11=R3-R5
; R1=R1+R3
; R2=R2+R3
; R3=R4-R1
; R4=R4+R1
; R1=R2-R5
; R1=R1+021
; R3=R3+R5
; R3=R3+021
; R1=R1+R2
; R4=R2-R1
; R1=R1+021
; Y11=(R3+R5)+021
; R2=R2+R4
; R2=R2+021
; Y11=(R4+R2)+021
; Y11=(R4+R2)+021
; Y11=(R4+R2)+021
; Y11=(R4+R2)+021
; LOOP BACK TO THE INNER LOOP
; INCREMENT REPEAT COUNTER FOR NEXT
; TIME
; IE=4+IE

```



# Appendix B2. fft\_4--Radix-4 Complex FFT to Be Called as a C Function

```

APPENDIX B2
NAME: fft_4 --- RADIX-4 COMPLEX FFT TO BE CALLED AS A C FUNCTION.
SYNOPSIS:
    int  fft_4(n, m, data)
    int  n      FFT SIZE: 2^M+4M
    int  m      NUMBER OF STAGES = LOG4(N)
    float *data  ARRAY WITH INPUT AND OUTPUT DATA
DESCRIPTION:
    GENERIC FUNCTION TO DO A RADIX-4 FFT COMPUTATION ON THE TRMS20C30.
    THE DATA ARRAY IS 2M+4M-LONG, WITH REAL AND IMAGINARY VALUES ALTERNATING.
    THE PROGRAM IS BASED ON THE FORTRAN PROGRAM IN THE BURRIS AND PARKS BOOK, P. 17.
    IN ORDER TO HAVE THE FINAL RESULT IN BIT-REVERSED ORDER, THE TWO
    MIDDLE BRANCHES OF THE RADIX-4 BUTTERFLY ARE INTERCHANGED DURING
    STORAGE. NOTE THIS DIFFERENCE WHEN COMPARING WITH THE PROGRAM ON
    P. 117. THE COMPUTATION IS DONE IN-PLACE, AND THE ORIGINAL DATA IS
    DESTROYED. BIT REVERSAL IS IMPLEMENTED AT THE END OF THE FUNCTION.
    IF THIS IS NOT NECESSARY, THIS PART CAN BE COODED OUT. THE
    SINE/COSINE TABLE FOR THE MIDDLE FACTORS IS EXPECTED TO BE SUPPLIED
    DURING LINK TIME, AND IT SHOULD HAVE THE FOLLOWING FORMAT:
        .global  _sine
        .data
        _sine  .float  value1 = sin(0*2pi/N)
              .float  value2 = sin(1*2pi/N)
              .....
              .float  value(SM/4) = sin((SM/4-1)*2pi/N)
    THE VALUES value1, value2, ETC., ARE THE SINE WAVE VALUES. FOR AN
    N-POINT FFT, THERE ARE NM/4 VALUES FOR A FULL AND A QUARTER PERIOD
    OF THE SINE WAVE. IN THIS WAY, A FULL SINE AND COSINE PERIOD ARE
    AVAILABLE (SUPERIMPOSED).
    STACK STRUCTURE UPON THE CALL:
        --FP(4) : DATA
        --FP(3) : M
        --FP(2) : N
        --FP(1) : RETURN ADDR
        --FP(0) : OLD FP
        ----->
    REGISTERS USED: R0, R1, R2, R4, R5, R6, R7, AR0, AR1, AR2, AR3, AR4,
    AR5, AR6, AR7, IRO, IRI, RS, RE, RC
    AUTHOR: PAMOS E. PAPPACHALIS
    TEXAS INSTRUMENTS
    OCTOBER 13, 1987
    *****

```

```

* FP          AR3
*
* .GLOBAL  _FFT_4
* .GLOBAL  _SINE
*
* .BSS    FFTSZ,1
* .BSS    LOFFT,1
* .BSS    INPUT,1
*
* .TEXT
*
* .SINTAB  .word  _SINE
*
* * INITIALIZE C FUNCTION
*
* _fft_4:  PUSH    FP
*          LDI    SP,FP
*          PUSH  R4
*          PUSH  R3
*          PUSH  R6
*          PUSH  R7
*          PUSH  AR4
*          PUSH  AR5
*          PUSH  AR6
*          PUSH  AR7
*
*          LDI    +FP(2),R0
*          STI
*          LDI    +FP(3),R0
*          STI
*          LDI    R0,LOFFT
*          LDI    +FP(4),R0
*          STI
*          R0,INPUT
*
*          * INITIALIZE FFT ROUTINE
*          .BSS  STAGE,1
*          .BSS  RPPOINT,1
*          .BSS  IEMIDX,1
*          .BSS  LPOINT,1
*          .BSS  JT,1
*          .BSS  IAI,1
*
*          LDI    @FFTSZ,R0
*          LDI    @FFTSZ,IRO
*          LDI    @FFTSZ,IRI
*          LDI    0,AR7
*          STI    AR7,@STAGE
*
*          LSH    1,IRO
*          LSH    -2,IRI
*          LDI    1,AR7
*          STI    AR7,@RPPOINT
*
*          * MOVE ARGUMENTS TO LOCATIONS MATCHING
*          * THE NAMES IN THE PROGRAM
*
*          * FFT STAGE #
*          * REPEAT COUNTER
*          * IE INDEX FOR SINE/COSINE
*          * SECOND-LOOP COUNT
*          * JT COUNTER IN PROGRAM, P. 117
*          * IAI INDEX IN PROGRAM, P. 117
*
*          * STAGE HOLDS THE CURRENT STAGE
*          * NUMBER
*          * IRO/IRI (BECAUSE OF REAL/IMAG)
*          * IRI=NI/4, POINTER FOR SINE/COS TABLE
*          * INITIALIZE REPEAT COUNTER OF FIRST
*          * LOOP

```

```

LSH      -2,RO
STI      AR7,@IENDX
ADDI     2,RO
STI      RO,@J
SUBI     2,RO
LSH      1,RO

*
*   OUTER LOOP
*
LOOP:
LDI      @INPUT,ARO
ADDI     RO,ARO,ARI
ADDI     R3,RI,R6
ADDI     RO,AR2,AR3
LDI      @POINT,RC
SUBI     1,RC

*
*   FIRST LOOP
*
RPTB    BLK1
ADDF    **AR0,**AR2,RI
ADDF    **AR2,**AR1,R3
ADDF    R3,RI,R6
SUBF    **AR2,**AR0,R4
STF     R6,**AR0
SUBF    R3,RI
SUBF    **AR1,R7
LDF     **AR1,R7
ADDF    **AR3,**AR1,R3
ADDF    R1,**AR1
ADDF    R5,**AR0,R2
SUBF    R3,RI,R6
SUBF    R6,**AR0,**(IRO)
SUBF    **AR3,**AR1,R6
SUBF    R7,**AR3,R3
SUBF    R1,**AR1,**(IRO)
SUBF    R6,RI,R5
STF     R5,**AR2
STF     R4,**AR3
SUBF    R3,R2,R5
ADDF    R3,R2
STF     R5,**AR2,**(IRO)
STF     R2,**AR3,**(IRO)

*
*   IF THIS IS THE LAST STAGE, YOU ARE DONE
*
LDI      @STAGE,AR7
ADDI     1,AR7
CPI     @LOFFT,AR7
BID     END
STI     AR7,@ISTAGE

```

```

*
*   MAIN INNER LOOP
*
LDI      1,AR7
STI      AR7,@IAI
LDI      2,AR7
STI      AR7,@LPONT

INLP:
LDI      2,AR6
ADDI     @LPONT,AR6
LDI      @LPONT,ARO
LDI      @IAI,AR7
ADDI     @IENDX,AR7
@INPUT,ARO
STI      AR7,@IAI
ADDI     RO,ARO,ARI
STI      AR6,@LPONT
ADDI     RO,AR1,AR2
ADDI     RO,AR2,AR3
LDI      @POINT,RC
SUBI     1,RC
CPI     @P1
BID     SPL
LDI      @IAI,AR7
LDI      @IAI,ARA
ADDI     @SINTAB,ARA
ADDI     AR4,AR7,AR5
SUBI     1,AR5
ADDI     AR7,AR5,AR6
SUBI     1,AR6

*
*   SECOND LOOP
*
RPTB    BLK2
ADDF    **AR2,**AR0,R3
ADDF    **AR3,**AR1,R5
ADDF    R5,R3,R6
SUBF    **AR2,**AR0,R4
SUBF    R5,R3
ADDF    **AR2,**AR0,RI
ADDF    **AR3,**AR1,R5
PPYF   R3,**AR3,(RI),R6
STF     R6,**AR0
ADDF    R5,RI,R7
SUBF    **AR2,**AR0,R2
SUBF    R5,RI
PPYF   R1,**AR5,R7
STF     R7,R6
SUBF    **AR3,**AR1,R5
PPYF   R1,**AR5,(RI),R7
STF     R6,**AR1
PPYF   R3,**AR5,R6
ADDF    R7,R6

```

```

; INIT IAI INDEX
; INIT LOOP COUNTER FOR INNER LOOP
; INCREMENT INNER LOOP COUNTER
; IAI=IAI+IE
; X(I),Y(I) POINTER
; X(I1),Y(I1) POINTER
; X(I2),Y(I2) POINTER
; X(I3),Y(I3) POINTER
; RC SHOULD BE ONE LESS THAN DESIRED #
; IF LPONT=J, GO TO
; SPECIAL BUTTERFLY
; CREATE COSINE INDEX ARA
; IAC=IAI+IAI-1
; IAC=IAC+IAI-1
; R2=Y(I1)+Y(I2)
; R5=Y(I1)+Y(I3)
; R6=R3+R5
; R4=Y(I1)-Y(I2)
; R3=R3-R5
; R1=X(I1)+X(I2)
; R2=X(I1)+X(I3)
; R6=R3+R5
; Y(I1)+R3+R5
; R7=R1+R5
; R2=X(I1)-Y(I2)
; R1=R1-R5
; X(I1)-R1+R5
; R6=R3+R2-R4+R5
; R5=Y(I1)-Y(I3)
; R7=R1+R2
; Y(I1)+R3+R2-R4+R5
; R6=R3+R5
; R6=R1-R2+R3+R5

```

```

; INITIALIZE IE INDEX
; JI=RO/2+2
; R0=RC
; ARO POINTS TO X(I)
; ARI POINTS TO X(I1)
; AR2 POINTS TO X(I2)
; AR3 POINTS TO X(I3)
; RC SHOULD BE ONE LESS THAN DESIRED #
; RI=Y(I1)+Y(I2)
; R3=Y(I1)+Y(I3)
; R6=R1+R3
; Y(I1)-Y(I2)
; Y(I1)+R1+R3
; R1=R1-R3
; R5=X(I2)
; R7=Y(I1)
; R3=X(I1)+X(I3)
; R1=X(I1)+X(I2)
; Y(I1)+R1-R3
; R6=R1+R3
; R2=X(I1)-X(I2)
; Y(I1)+R1+R3
; R1=R1-R3
; R6=X(I1)-X(I3)
; R5=R4+R6
; Y(I2)+R4-R6
; Y(I3)+R4+R6
; R5=R2-R3
; R2=R2+R3
; X(I2)+R2+R3
; X(I3)+R2+R3

```

```

; CURRENT FFT STAGE

```



POP AR7  
POP AR6  
POP AR5  
POP AR4  
POP R7  
POPF R6  
POP R5  
POP R4  
POP FP  
RETS

## **Appendix C.Radix-2 Real FFT**

# Appendix C1. Generic Program to Do a Radix-2 Real FFT Computation on the TMS320C30

```

*
* APPENDIX C1
*
* GENERIC PROGRAM TO DO A RADIX-2 REAL FFT COMPUTATION ON THE TMS320C30
*
* THE PROGRAM IS TAKEN FROM THE PAPER BY SØRENSEN ET AL., JUNE 1987 ISSUE
* OF THE TRANSACTIONS ON ASSP.
*
* THE (REAL) DATA RESIDE IN INTERNAL MEMORY. THE COMPUTATION IS DONE
* IN-PLACE. THE BIT REVERSAL IS DONE AT THE BEGINNING OF THE PROGRAM.
*
* THE TRUNDLE FACTORS ARE SUPPLIED IN A TABLE PUT IN A DATA SECTION. THIS
* DATA IS INCLUDED IN A SEPARATE FILE TO PRESERVE THE GENERIC NATURE OF THE
* PROGRAM. FOR THE SAME PURPOSE, THE SIZE OF THE FFT (N AND LOG2(N)) ARE
* DEFINED IN A .GLOBAL DIRECTIVE AND SPECIFIED DURING LINKING. THE LENGTH OF
* THE TABLE IS N/4 + N/4 = N/2.
*
* AUTHOR: PANOS E. PAPANICOLAIS
* TEXAS INSTRUMENTS
* SEPTEMBER 8, 1987
*
* .GLOBAL FFT
* .GLOBAL N
* .GLOBAL SINE
* .GLOBAL SINE
*
* .USECT "IN",1024
* .BSS OUTP,1024
*
* .TEXT
*
* INITIALIZE
*
* .WORD FFT
*
* .SPACE 100
*
* FFTSIZ .WORD N
* LOGFFT .WORD M
* SINTAB .WORD SINE
* INPUT .WORD INP
* OUTPUT .WORD OUTP
*
* FFT: LIP FFTSIZ
*
* DO THE BIT-REVERSING AT THE BEGINNING
*
* LDI @FFTSIZ,RC
* SUBI 1,RC
* LSH -1,IRO
* LDI @INPUT,ARO
* LDI @INPUT,ARI
*
* RPTB BITRV
*
* ; CHANGE LOCATIONS ONLY
* ; IF ARO<ARI
*
* ARI,ARO
* CONT
* **ARO,RO
* **ARI,RI
* RO,**ARI
* RI,**ARO
* **ARO,**IRO,B
*
* LENGTH-TWO BUTTERFLIES
*
* LDI @INPUT,ARO
* LDI IRO,RC
* SUBI 1,RC
*
* RPTB BLK1
*
* ADFI **ARO,**ARO,**RO
* SUBF **ARO,**ARO,RI
* STF RO,**ARO
* STF RI,**ARO**
*
* BLK1
*
* ; FIRST PASS OF THE DD-20 LOOP (STAGE K=2 IN DD-10 LOOP)
*
* LDI @INPUT,ARO
* LDI 2,IRO
* LDI @FFTSIZ,RC
* LSH -2,RC
* SUBI 1,RC
*
* RPTB BLK2
*
* ADFI **ARO(IRO),**ARO**(IRO),RO ; RO=X(1)+X(1+2)
* SUBF **ARO,**ARO,(IRO),RI ; RI=X(1)-X(1+2)
* NEGF **ARO,RO ; RO=-X(1+3)
* STF RO,**ARO(IRO) ; X(1)=X(1)+X(1+2)
* STF RI,**ARO**(IRO) ; X(1+2)=X(1)-X(1+2)
* STF RO,**ARO ; X(1+3)=X(1+3)
*
* MAIN LOOP (FFT STAGES)
*
* LDI @FFTSIZ,IRO
* LSH -2,IRO
* LDI 3,RS
* LDI 1,RA
* LDI 1,RA
* LDI 2,R3
* LSH -1,IRO
* LSH 1,RA
* LSH 1,R3
*
* LOOP
*
* INNER LOOP (DD-20 LOOP IN THE PROGRAM)
*
* LDI @INPUT,AS
* LDI IRO,ARO
* ADI @SINTAB,ARO
* LDI RA,IRI
*
* ; XCHANGE LOCATIONS ONLY
* ; IF ARO<ARI
*
* ARI,ARO
* CONT
* **ARO,RO
* **ARI,RI
* RO,**ARI
* RI,**ARO
* **ARO,**IRO,B
*
* LENGTH-TWO BUTTERFLIES
*
* LDI @INPUT,ARO
* LDI IRO,RC
* SUBI 1,RC
*
* RPTB BLK1
*
* ADFI **ARO,**ARO,**RO
* SUBF **ARO,**ARO,RI
* STF RO,**ARO
* STF RI,**ARO**
*
* BLK1
*
* ; FIRST PASS OF THE DD-20 LOOP (STAGE K=2 IN DD-10 LOOP)
*
* LDI @INPUT,ARO
* LDI 2,IRO
* LDI @FFTSIZ,RC
* LSH -2,RC
* SUBI 1,RC
*
* RPTB BLK2
*
* ADFI **ARO(IRO),**ARO**(IRO),RO ; RO=X(1)+X(1+2)
* SUBF **ARO,**ARO,(IRO),RI ; RI=X(1)-X(1+2)
* NEGF **ARO,RO ; RO=-X(1+3)
* STF RO,**ARO(IRO) ; X(1)=X(1)+X(1+2)
* STF RI,**ARO**(IRO) ; X(1+2)=X(1)-X(1+2)
* STF RO,**ARO ; X(1+3)=X(1+3)
*
* MAIN LOOP (FFT STAGES)
*
* LDI @FFTSIZ,IRO
* LSH -2,IRO
* LDI 3,RS
* LDI 1,RA
* LDI 1,RA
* LDI 2,R3
* LSH -1,IRO
* LSH 1,RA
* LSH 1,R3
*
* LOOP
*
* INNER LOOP (DD-20 LOOP IN THE PROGRAM)
*
* LDI @INPUT,AS
* LDI IRO,ARO
* ADI @SINTAB,ARO
* LDI RA,IRI

```

```

*
LDI    AR5,AR1
ADDI   L,AR1
LDI    R3,AR3
LDI    R3,AR3
SUBI   R3,AR3,AR4
LDI    R3,AR2,AR4
LUF    +AR5+(IR1),R0
      +AR5(IR1),R0,R1
      R0,++AR5(IR1),R0
      R1,++AR5(IR1)
      R1,X(1)+X(1+M2)
      X(1)+X(1)+X(1+M2)
      R0=X(1)-X(1+M2)
      ++AR5(IR1),R1
      R0,AR5
      R1,AR5
*      X(1+M+M2)=-(1+M+M2)
*
* INNERMOST LOOP
LDI    @FFTSIZ,IR1
LSH   -2,IR1
LDI    R4,RC
SUBI   R4,RC
*      IR1=SEPARATION BETWEEN SIN/COS TABLS
*      REPEAT NM-1 TIMES
*
RPTB   BLK3,++AR6(IR1),R0
RYF   ++AR4,AR6,R1
RYF   ++AR4,++AR6(IR1),R1
AUF   R0,RI,R2
RYF   ++AR3,++AR6+(IR0),R0
SUBF  R0,RI,R0
AUF   ++AR2,R0,R1
AUF   ++AR2,R0,R1
AUF   ++AR3++
STF   R1,AR63++
STF   R2,AR61,RI
STF   R1,AR61++
STF   R1,AR62--
*      R1=X(13)*COS
*      R1=X(14)*SIN
*      R2=X(13)*COS+X(14)*SIN
*      R0=X(13)*SIN
*      R0=-X(13)*SIN+X(14)*COS
*      R1=X(12)*AR0
*      R1=X(11)*AR2
*      R1=X(12)*AR0
*      R2=X(11)*R2
*      R1,AR61++
*      X(12)=X(11)*R2
*
SUBI   @INPUT,AR5
ADDI   R4,AR5
CPI    @FFTSIZ,AR5
BLT    BLTO
LDI    @INPUT,AR5
NOP
NOP
*
*
ADDI   L,AS
CPI    @OFFSET,AS
BLE   LOOP
NOP
NOP
*

```

```

*
NOP
NOP
*
END
BR    END
; BRANCH TO ITSELF AT THE END
;
NOP
NOP

```

# Appendix C2. fft\_rl—Radix-2 Real FFT to Be Called as a C Function

```

APPENDIX C2
NAME:
    fft_rl ---- RADIX-2 REAL FFT TO BE CALLED AS A C FUNCTION.
SYNOPSIS:
    int fft_rl(int n, data)
    int n      FFT SIZE: N/2+1
    int m      NUMBER OF STAGES = LOG2(N)
    float data  ARRAY WITH INPUT AND OUTPUT DATA
DESCRIPTION:
    GENERIC FUNCTION TO DO A RADIX-2 FFT COMPUTATION ON THE TMS320C20.
    THE DATA ARRAY IS M-LONG, WITH ONLY REAL DATA. THE OUTPUT IS STORED
    IN THE SAME LOCATIONS WITH REAL AND IMAGINARY POINTS R AND I AS
    FOLLOWS: R(0), R(1), ..., R(N/2), I(N/2+1), ..., I(1)
    THE PROGRAM IS BASED ON THE FORTRAN PROGRAM IN THE PAPER BY SORBERSEN
    ET AL., JUNE 1987 ISSUE OF TRANS. ON ASSP. THE COMPUTATION IS DONE
    IN-PLACE, AND THE ORIGINAL DATA IS DESTROYED. BIT REVERSAL IS
    IMPLEMENTED AT THE BEGINNING OF THE FUNCTION. IF THIS IS NOT
    NECESSARY, THIS PART CAN BE COMMENTED OUT.
    THE SINE/COSINE TABLE FOR THE WHOLE FACTORS IS EXPECTED TO BE
    SUPPLIED DURING LINK TIME, AND IT SHOULD HAVE THE FOLLOWING FORMAT:
        .global  _sine
        .data
        .float value1 = sin(0+2*pi/N)
        .....
        .float value(N/2) = cos((N/4)+2*pi/N)
    THE VALUES value1 TO value(N/4) ARE THE FIRST QUARTER OF THE SINE
    PERIOD AND value(N/4+1) TO value(N/2) ARE THE FIRST QUARTER OF THE
    COSINE PERIOD.
    STACK STRUCTURE UPON THE CALL:
        +-----+
        -FP(4) : DATA
        -FP(3) : N
        -FP(2) : N
        -FP(1) : RETURN ADDR
        -FP(0) : OLD FP
        +-----+
REGISTERS USED: R0, R1, R2, R3, R4, R5, AR0, AR1, AR2, AR4, AR5, IRO,
                IRI, RS, RE, RC
AUTHOR: PANOS E. PAPAIOACHALIS
        TEXAS INSTRUMENTS
OCTOBER 13, 1987
*****

```

```

* PP
* .SET AR3
* .GLOBAL _FFT_BL
* .GLOBAL _SINE
* : ENTRY POINT FOR EXECUTION
* : ADDRESS OF SINE TABLE
*
* .BSS FFTSIZ,1
* .BSS LOFFFT,1
* .BSS INPUT,1
* .TEXT
*
* SINTAB .word _SINE
*
* INITIALIZE C FUNCTION
*
* _FFT_BL: PUSH PP
*         LODI SP,PP
*         PUSH R4
*         PUSH R5
*         PUSH AR4
*         PUSH AR5
*
*         LODI 4-PP(2),R0
*         STI R0,FFTISZ
*         LODI 4-PP(3),R0
*         STI R0,LOFFFT
*         LODI 4-PP(4),R0
*         STI R0,INPUT
*
* DO THE BIT REVERSING AT THE BEGINNING
*
*         LODI 4-FFTISZ,RC
*         SUBI 1,RC
*         LODI 4-FFTISZ,IRO
*         LSH -1,IRO
*         LODI 4-INPUT,ARO
*         LODI 4-INPUT,ARI
*
*         RPTB BITRV
*         CPTI ARI,ARO
*         BBE CONT
*         LUF *ARO,R0
*         LUF *ARI,R1
*         STF R0,*ARI
*         STF R1,*ARO
*         NOP *ARO++
*         NOP *ARI++(IRO)B
*
* LENGTH-TWO BUTTERFLIES
*
*         LODI 4-INPUT,ARO
*         LODI 4-IRO,RC
*         SUBI 1,RC

```



```

*
PFTB BLK1
ADEF ++ARO, ++ARO++, RO
SUBF ++RO, ++ARO, RI
SIF RO, ++ARO
SIF RI, ++ARO++
BLK1
::
*
* FIRST PASS OF THE DO-20 LOOP (STAGE K=2 IN DO-10 LOOP)
*
LDI @INPUT, ARO
LDA 2, IRO
LDI @FFTSIZ, IC
LSH -2, IC
SUBI 1, IC
*
* ARO POINTS TO X(1)
* IRO=2*IC
* REPEAT N/4 TIMES
* IC SHOULD BE ONE LESS THAN DESIRED #
*
BLK2
PFTB ++ARO(IRO), ++ARO++(IRO), RO ; RO=(I(1)+(I+2))
ADEF ++RO, ++ARO(IRO), RI ; RI=X(I)-X(I+2)
SUBF ++ARO, RO
MEBF ++ARO, RO ; RO=-X(I+2)
SIF RO, ++ARO(IRO) ; X(I)=X(I)+X(I+2)
SIF RI, ++ARO++(IRO) ; X(I+2)=X(I)-X(I+2)
SIF RO, ++ARO ; X(I+2)=-X(I+2)
*
* MAIN LOOP (FT STAGES)
*
LDI @FFTSIZ, IRO
LSH -2, IRO
LDI 3, IS
LDI 1, PA
LDI 1, RA
LDI 2, R3
LSH -1, IRO ; E=E/2
LSH 1, PA ; M=2*MA
LSH 1, R3 ; N2=2*MC
*
* INNER LOOP (DO-20 LOOP IN THE PROGRAM)
*
LDI @INPUT, ARS
LDI IRO, ARO
AADI @ESIMUL, ARO ; ARO POINTS TO SIM/COS TABLE
LDI RA, IRI
*
LDI ARS, AR1
AADI 1, AR1 ; AR1 POINTS TO X(11)=(I+J)
LDI AR1, AR3
AADI R3, AR3 ; AR3 POINTS TO X(13)=(I+J*MC)
LDI AR3, AR2
SUBI 2, AR2 ; AR2 POINTS TO X(12)=X(I-J*MC)
AADI R3, AR2, AR4 ; AR4 POINTS TO X(14)=X(I-J*MI)
*
LDF ++ARS+(IR1), RO ; RO=X(1)
ADEF ++ARS(IR1), RO, RI ; RI=X(1)+X(1+MC)
SUBF RO, ++ARS(IR1), RO ; RO=++ARS(IR1), RO
SIF RI, ++ARS(IR1) ; X(1)=X(1)+X(1+MC)
MEBF RO
*
::
*
* SEPARATION BETWEEN SIN/COS TABS
*
LDI @FFTSIZ, IRI
LSH -2, IRI
LDI RA, IC
SUBI 2, IC
*
* REPEAT M-1 TIMES
*
BLK3
PFTB ++R3, ++ARO(IR1), RO ; RO=X(13)+COS
PFTB ++RA, ++ARO, RI ; RI=X(14)+SIN
PFTB ++RA, ++ARO(IR1), RI ; RI=X(14)+COS
ADEF RO, RI, R2 ; R2=X(13)+COS*X(14)+SIN
SUBF ++R3, ++ARO++(IRO), RO ; RO=X(13)+SIN
SUBF ++R2, RO, RI ; RO=-X(13)+SIN*X(14)+COS ...
SUBF ++R2, RO, RI ; RI=X(12)+RO ...
SIF RI, ++R3++ ; X(13)=-(I2)+RO ...
SIF RI, ++R4-- ; RI=X(12)+RO ...
SIF R2, AR1, RI ; X(14)=X(12)+RO ...
SIF RI, AR1++ ; X(11)=X(11)+R2
SIF RI, AR2-- ; X(12)=X(11)-R2
BLK3
*
* INPUT, ARS
* ARO=I+MI
* LOOP BACK TO THE INNER LOOP
*
LDI @INPUT, ARS
AADI R4, AR5
CPY @FFTSIZ, ARS
BLTD INLUP
AADI @INPUT, ARS
MCP
*
* RESTORE THE REGISTER VALUES AND RETURN
*
AADI 1, RS
CPY @GLOFFT, IC
BLE LOOP
*
* RESTORE THE REGISTER VALUES AND RETURN
*
POP ARS
POP AR4
POP RS
POP R4
POP PP
PP
RETS

```

# Appendix C3. Generic Program to Do a Radix-2 Real Inverse FFT Computation on the TMS320C30

```

*
* APPENDIX C3
*
* GENERIC PROGRAM TO DO A RADIX-2 REAL INVERSE FFT COMPUTATION ON THE
* TMS320C30.
*
* THE REAL DATA RESIDE IN INTERNAL MEMORY. THE COMPUTATION IS DONE
* IN-PLACE. THE BIT REVERSAL IS DONE AT THE BEGINNING OF THE PROGRAM. THE
* INPUT DATA ARE STORED IN THE FOLLOWING ORDER:
*
* RE(0), RE(1),..., RE(N/2), IM(N/2-1),..., IM(1)
*
* THE TWIDDLE FACTORS ARE SUPPLIED IN A TABLE PUT IN A .DATA SECTION. THIS
* DATA IS INCLUDED IN A SEPARATE FILE TO PRESERVE THE GENERIC NATURE OF THE
* PROGRAM. FOR THE SAME PURPOSE, THE SIZE OF THE FFT N AND LOG2(N) ARE
* DEFINED IN A .GLOBAL DIRECTIVE AND SPECIFIED DURING LINKING. THE LENGTH OF
* THE TABLE IS N/4 * N/4 = N/2.
*
* AUTHOR: PANGS PAPAETHALIS
* TEXAS INSTRUMENTS
* DECEMBER 21, 1988
*
* .GLOBAL IFFT
* .GLOBAL N
* .GLOBAL M
* .GLOBAL LOG2(N)
* .GLOBAL SINE
*
* .BSS TMP,1024
*
* .TEXT
*
* INITIALIZE
*
* .WORD IFFT
*
* .SPACE 100
*
* .WORD N
* .WORD LOG2(N)
* .WORD SINE
* .INPUT .WORD TMP
*
* IFFT: LOP FFTSIZ
*
* MAIN LOOP (FFT STAGES)
*
* LDI 1,IRO
* LDI 3,RS
* LDI @FFTSIZ,R3
* LSH -1,R3
* LDI @FFTSIZ,IM
* LSH -2,R4
*
* INNER LOOP
*
* @INPUT,ARS
* IRO,ARO
* @SINTAB,ARO
* R4,IR1
*
* ARS POINTS TO X(1)
* ARO POINTS TO SIN/COS TABLE
* IRI=MM
*
* AR1 POINTS TO X(11)=X(1+N)
*
* AR3 POINTS TO X(13)=X(1+N*H2)
*
* AR2 POINTS TO X(12)=X(1-1*H2)
*
* AR4 POINTS TO X(14)=X(1-1*H4)
*
* * POINT TO X(1+MM)
* ++ARS(IR1)
* ++ARS(IR1),RO
* ++ARS(IR1),@ARS(IR1),RI
* RO,@ARS(IR1)
* STP RI,++ARS(IR1)
* LIF ++ARS(IR1)
* PIVF 2.0,RO
* STP RI,@ARS(IR1)
* LIF ++ARS(IR1)
* PIVF -2.0,RI
*
* X(1+MM)=2*X(1+MM)
*
* X(1+MM*H2)=X(1+MM*H2)*2
*
* * INNERMOST LOOP
*
* LDI @FFTSIZ,IR1
* LSH -2,IR1
* LDI R4,RC
* SUBI 2,RC
*
* IRI=SEPARATION BETWEEN SIN/COS TABLS
* REPEAT MM-1 TIMES
*
* RPTB
* SUBF #R2,*AR1,R1
* ADF R1,*AR0,(R1),RO
* PIVF R0,*AR1++
* ADF #AR3,*AR4,R2
* SUBF #AR3,*AR4,R6
* PIVF R2,*AR0,R6
* STP R6,*AR2--
* SUBF R6,RO
* PIVF R2,*AR0,(R1),R6
* STP R1,*AR3++
* PIVF R1,*AR0++(IRO),RO
* ADF R6,RO
* STP R0,*AR4--
*
* R6=T2*SIN
* X(12)=X(14)-(13)
*
* R6=T2*COS
* X(13)=T1*COS-T2*SIN
*
* R6=T1*SIN
* X(14)=T1*SIN+T2*COS
*
* * LOOP BACK TO THE INNER LOOP
*
* @INPUT,ARS
* IRO,ARO
* @SINTAB,ARO
*
* ARO POINTS TO SIN/COS TABLE

```

; IF #RC#GT#R1

CONT #ARO,R0  
LDF #ARI,R1  
R0,#R1  
STF R1,#ARO  
R1,#ARO  
CONT #ARO++  
BITRV #ARI++(R0)B  
\*  
END

; BRANCH TO ITSELF AT THE END

BR .END

.END

ADDI 1,R5  
CMP1 @C(OEFTT),R5  
BLED LOOP  
LSH 1,1R0  
LSH -1,R4  
LSH -1,R3

LAST PASS OF THE MAIN LOOP

LDI @INPUT,ARO  
LDI 2,1R0  
LDI @FFTSIZ,RC  
LSH -2,RC  
SUBI 1,RC

RC SHOULD BE ONE LESS THAN DESIRED #

RO=X(1+2)

LDF #ARO(1R0),RO  
RPTB BLK2  
ADDF RO,#ARO++(1R0),R1  
SUBF RO,#ARO(1R0),R1  
STF R1,#ARO(1R0)  
STF R1,#ARO++  
LDF #ARO,R1  
PPYF 2,0,R1  
STF R1,#ARO(1R0)  
LDF #ARO++,R1  
PPYF -2,0,R1  
STF R1,#ARO  
LDF #ARO(1R0),RO

RI=X(1)+X(1+2)  
RI=X(1)-X(1+2)  
X(1)=X(1)+X(1+2)  
X(1+2)=X(1)-X(1+2)  
R1=2.0#X(1+1)  
X(1+1)=2.0#X(1+1)  
R1=-2.0#X(1+3)  
X(1+3)=-2.0#X(1+3)  
R0=X(1+4#2)

LENGTH-TWO BUTTERFLIES

LDI @INPUT,ARO  
LDI @FFTSIZ,RC  
LSH -1,RC  
SUBI 1,RC

RC SHOULD BE ONE LESS THAN DESIRED #

RPTB BLK1  
ADDF #ARO,#ARO++,RO  
SUBF #ARO,#ARO,R1  
STF R1,#ARO  
STF R1,#ARO++

DO THE BIT REVERSING AT THE END

LDI @FFTSIZ,RC  
SUBI 1,RC  
LDI @FFTSIZ,1R0  
LSH -1,1R0  
LDI @INPUT,ARO  
LDI @INPUT,ARI

RC=M  
RC SHOULD BE ONE LESS THAN DESIRED #  
IRO=HALF THE SIZE OF FFT=M/2

RPTB BITRV  
ARI,ARI

CHANGE LOCATIONS ONLY

## **Appendix D. Discrete Hartley Transform**

# Appendix D1. Generic Program to Do a Radix-2 Hartley Transform on the TMS320C30

```

* APPENDIX D1
*
* GENERIC PROGRAM TO DO A RADIX-2 HARTLEY TRANSFORM ON THE TMS320C30.
*
* THE PROGRAM IS TAKEN FROM THE PAPER BY SØRENSEN ET AL., OCT 1985 ISSUE
* OF THE TRANSACTIONS ON ASSP.
*
* THE (REAL) DATA RESIDE IN INTERNAL MEMORY. THE COMPUTATION IS DONE
* IN-PLACE. THE BIT-REVERSAL IS DONE AT THE BEGINNING OF THE PROGRAM.
*
* THE TWIDDLE FACTORS ARE SUPPLIED IN A TABLE PUT IN A .DATA SECTION. THIS
* DATA IS INCLUDED IN A SEPARATE FILE TO PRESERVE THE GENERIC NATURE OF THE
* PROGRAM. FOR THE SAME PURPOSE, THE SIZE OF THE FFT N AND LOG2(N) ARE
* DEFINED IN A .GLOBL DIRECTIVE AND SPECIFIED DURING LINKING. THE LENGTH OF
* THE TABLE IS N/4 + N/4 = N/2.
*
* AUTHOR: PANOS PAPAIOACHALIS           DECEMBER 14, 1988
* TEXAS INSTRUMENTS
*
* .GLOBL  FRT          ; ENTRY POINT FOR EXECUTION
* .GLOBL  N            ; FFT SIZE
* .GLOBL  M            ; LOG2(N)
* .GLOBL  SINE         ; ADDRESS OF SINE TABLE
*
* .BSS   IMP, LOG2    ; MEMORY WITH INPUT DATA
*
* .TEXT
*
* INITIALIZE
*
* .WORD  FRT          ; STARTING LOCATION OF THE PROGRAM
*
* .SPACE 100         ; RESERVE 100 WORDS FOR VECTORS, ETC.
*
* .WORD  N            ; WORD N
* LOGFHT .WORD  M      ; WORD M
* SINTAB .WORD  SINE   ; WORD SINE
* INPUT  .WORD  IMP     ; WORD IMP
*
* FRT:  LDP  FRTSIZ, RC ; COMMAND TO LOAD DATA PAGE POINTER
*
* DO THE BIT REVERSING AT THE BEGINNING
*
*   LDI  MPTSIZ, RC   ; RC=N
*   SUBI 1, RC       ; RC SHOULD BE ONE LESS THAN DESIRED #
*   LSH  -1, IR0     ; IR0=HALF THE SIZE OF FRT=N/2
*   LDI  IRINPUT, ARO ; INPUT ARO
*   LDI  RINPUT, ARI ; INPUT ARI
*
* BITRV BITRV          ; BITRV LOCATIONS ONLY
*   ARI, ARO          ; IF ARO<ARI
*   COMT             ;
*   BBE             ;
*   LIF             ;
*   ARO, RO          ;
*
* .GLOBL  ARI, ARI    ; ARO POINTS TO X(1)
* .GLOBL  RO, ARI    ;
* .GLOBL  RI, ARO    ;
* .GLOBL  ARO++, RO  ;
*
* BITRV BITRV          ; LENGTH-TWO BUTTERFLIES
*   ARI++, IR0B      ;
*
*   LDI  RINPUT, ARO ; ARO POINTS TO X(1)
*   LDI  IR0, RC     ; REPEAT N/2 TIMES
*   SUBI 1, RC      ; RC SHOULD BE ONE LESS THAN DESIRED #
*
*   RPTB ++ARO, ARO++, RO ; RO=X(1)+X(1+1)
*   ADFB ARO, ARO, RI ; RI=X(1)-X(1+1)
*   SUBF ARO, ARO, RI ; RI=X(1)+X(1+1)
*   STF  RO, ARO    ; X(1)+X(1)+X(1+1)
*   STF  RI, ARO++  ; X(1)+X(1)-X(1+1)
*
*   FIRST PASS OF THE 1D-30 LOOP (STAGE K=2 IN 1D-20 LOOP)
*
*   LDI  RINPUT, ARO ; ARO POINTS TO X(J)
*   LDI  2, IR0      ; IR0=2*MC
*   LDI  MPTSIZ, RC  ;
*   LSH  -2, RC      ; REPEAT N/4 TIMES
*   SUBI 1, RC       ; RC SHOULD BE ONE LESS THAN DESIRED #
*
*   BLK2 RPTB        ;
*   ADFB ++ARO(IR0), ARO+++(IR0), RO ; RO=X(J)+X(L2)
*   SUBF ARO, ARO, ARO(IR0), RI ; RI=X(J)-X(L2)
*   STF  ++ARO, RO   ; RO=X(L)
*   ADFB ++ARO, RO   ; RO=X(L)
*   ADFB RO, ARO, RI ; RI=X(L3)+X(L4)
*   SUBF RI, ARO++   ; X(L2)+X(J)-X(L2)
*   STF  RO, ARO(IR0), RI ; RI=X(L3)-X(L4)
*   STF  RI, ARO++   ; X(L3)+X(L3)+X(L4)
*   BLK2 RPTB        ; X(L4)+X(L3)-X(L4)
*
*   MAIN LOOP (FRT STAGES)
*
*   LDI  MPTSIZ, IR0 ; IR0=INDEX FOR E
*   LSH  -2, IR0    ;
*   LDI  3, RC      ; RC HOLDS THE CURRENT STAGE NUMBER
*   LDI  1, RA      ; RA=M
*   LDI  2, R3      ; R3=MC
*   LDI  2, R3      ; R3=MC
*   LSH  -1, IR0   ; E=E/2
*   LSH  1, RA     ; RA=2*RA
*   LSH  1, R3     ; R3=2*R3
*
*   TIMER LOOP (1D-30 LOOP IN THE PROGRAM)
*
*   LDI  RINPUT, ARI ; ARI POINTS TO X(J)
*   LDI  IR0, ARO   ; ARO POINTS TO SIN/COS TABLE
*   ADFI RINTAB, ARO ;
*   LDI  RA, IRI    ; IRI=M

```



## Appendix E. Discrete Cosine Transform

# Appendix E1. A Fast Cosine Transform

```

* APPENDIX E1
*
* A FAST COSINE TRANSFORM
*
* BASED ON THE ALGORITHM OUTLINED BY BYEONG GI LEE IN HIS ARTICLE, FCT - A
* FAST COSINE TRANSFORM, PUBLISHED IN THE PROCEEDINGS OF THE IEEE INTER-
* NATIONAL CONFERENCE ON ACOUSTICS, SPEECH, AND SIGNAL PROCESSING, SAN
* DIEGO, CA, 19-21 MARCH 1984, P. 284-3/1-4 VOL. 2, (CH1954-5/84/0000-0299),
*
* LEE'S ALGORITHM HAS BEEN MODIFIED TO ALLOW NATURAL ORDER TIME DOMAIN
* COEFFICIENTS RATHER THAN THE LESS ORDERED INPUT SUGGESTED IN HIS ARTICLE.
*
* THE FREQUENCY DOMAIN COEFFICIENTS ARE IN BIT REVERSE ORDER. THIS IS AN IN
* PLACE CALCULATION.
*
* AUTHOR: PAUL WILHELM
*
*       .globl FCT
*       : FAST COSINE TRANSFORM ENTRY POINT.
*       .globl M
*       : LENGTH OF DATA ENTRY.
*       .globl COS_TAB
*       : TABLE OF COSINE COEFFICIENTS.
*       .globl COEFF
*       : TABLE OF INPUT DATA.
*
*       .text
*
* FCTSIZE word M
* _COS_   word COS_TAB
* _DATA_  word COEFF
*
* FCT:
* LD1     @FCTSIZE, #R0
* LD1     @FCTSIZE, #R1
* LD1     @DATA, #R6
* LD1     @COS, #R7
* LD1     @R0, #R1
* LD1     @-1, #R0
* LD1     @R6, #R1
* LD1     @R6, #R0, #R2
* SUB1    1, #R2
* LSH3    1R0, #R0, #R3
* LD1     1, #R5
* ADD1    @R6, #R2
* ADD13   1R0, #R3, #R4
* ADD13   1R0, #R5, #R1C
*
* FIRST LOOP SERIES
*
* THIS LOOP SERIES DOES ALL THE BUTTERFLY STAGES EXCEPT THE FINAL ONE.
*
* RPT8    END_CENTER_LOOP
  
```

```

* OUTSIDE_LOOP:
* MIDDLE_LOOP:
*
*
* @R2, #R2
* @R3, #R3
*
* SUBF3   @R3, @R4, #R1
* SUBF3   @R2, @R1, #R0
* SUBTRACT SECOND BUTTERFLY DATA.
* SUBTRACT FIRST BUTTERFLY DATA.
*
* MPYF3   R1, #+*#R7, #R1
* MPYF3   R2, @R4, #R3
* MULTIPLY 2ND SUBTRACTION RESULT BY
* COSINE COEFFICIENT. ADD SECOND
* BUTTERFLY DATA.
*
* MPYF3   R0, #+*#R7, #R0
* ADDF3   R2, @R4, #R2
* MULTIPLY 1ST SUBTRACTION RESULT BY
* COSINE COEFFICIENT. ADD FIRST
* BUTTERFLY DATA.
*
* STF     R1, @R2+*(1R1)/2
* STF     R3, @R4+*(1R1)/2
* SAVE 2ND MULTIPLY RESULT IN LOWER
* HALF IF BUTTERFLY. SAME 2ND
* ADDITION IN UPPER 2ND BUTTERFLY.
*
* END_CENTER_LOOP:
*
* STF     R0, @R3+*(1R1)/2
* STF     R2, @R1+*(1R1)/2
* SAVE 1ST MULTIPLY IN LOWER HALF OF
* 2ND BUTTERFLY. SAME 1ST ADDITION
* IN UPPER 1ST BUTTERFLY.
*
* END OF CENTER LOOP OF FIRST LOOP SERIES.
*
* ADD13   1R0, #R5, #R1C
*
* ADDF3   @R3+*, @R2--, #R0
* CPY1    @R3, #R2
* UPDTE DATA POINTERS.
* HAVE BUTTERFLIES BEEN COMPLETED?
*
* BR1D    MIDDLE_LOOP
* DELAYED BRANCH, IF NOT.
*
* ADDF3   @R1+*, @R4--*, #R0
* UPDTE FINAL TWO POINTERS FOR NEXT
* REPEAT.
*
* AND1    2, #R7
* OR      01000, #T1
* UPDTE COSINE COEFFICIENT POINTER.
* SET REPEAT MODE. (FASTER THAN USING
* RPT8 WHEN START AND END ADDRESS
* ARE STILL GOOD)
*
* DELAY BRANCH FROM HERE TO MIDDLE_LOOP.
*
* LSH     -1, #R1
* LD1     @R6, #R1
* UPDTE INDEX REGISTER. (DIVIDE BY 2)
* REINITIALIZE DATA POINTERS.
*
* AND1    1R0, #R4, #R2
* AND1    1R1, #R2
*
* CPY1    2, #R1
* BR1D    OUTSIDE_LOOP
* IS FIRST BUTTERFLY SERIES COMPLETE?
*
* LSH     1, #R5
* SUB13   1R0, #R4, #R3
* MULTIPLY 2'S POWER COUNTER BY 2.
* CONTINUE REINITIALIZING DATA
* POINTERS.
*
* ADD13   1R0, #R5, #R1C
* SET REPEAT COUNTER FOR REPEAT BLOCK.
*
* END OF FIRST LOOP SERIES.
*
* FINAL BUTTERFLY STAGE_LOOP.
  
```



```

* INCLUDES LAST BUTTERFLIES AND FIRST STAGE OF BIT REVERSE ADDITIONS.
*
*      4,R1           ; INITIALIZE INDEX REGISTER.
*      LDI 1,AR3     ; SET UP DATA POINTERS.
*      LSH -1,AR5
*      ADDI 3,AR4
*      ADDI 100,AR5,RC
*      MPPF3 #487,++R7,R4
*      ; INITIALIZE REPEAT COUNTER,
*      ; CALCULATE (207/405)(P/4),
*      ; (I.E.-> (SR)(231)/M THIS VALUE IS
*      ; CALLED, S, BELOW.)
*      RPTB END_2ND_LOOP
*      ; TWO BUTTERFLIES ARE CALCULATED PER
*      ; LOOP.
*
*      SUBF3 #482,++R1,R0
*      SUBF3 #484,++R2,R1
*      MPPF3 R0,R4,R0
*      ADDF3 #483++(R1),++R4++(R1),R3 ; BY S, AND 2ND BUTTERFLY
*      ; DATA.
*      MPPF3 R1,R4,R1
*      ADDF3 #481++(R1),++R2++(R1),R2 ; BY S, AND 1ST BUTTERFLY
*      ; DATA.
*      MPPF3 R3,++R7,R3
*      STF R0,++R2(R1)
*      ; 7071, SAME 1ST SUBTRACTION IN
*      ; LOWER 1/2 OF 1ST BUTTERFLY.
*      MPPF3 R2,++R7,R2
*      STF R1,++R4(R1)
*      ; 7071, SAME 2ND SUBTRACTION IN
*      ; LOWER 1/2 OF 2ND BUTTERFLY.
*      ADDF3 R3,R1,R3
*      ; AND 2ND SUBTRACTION MULTIPLY TO 2ND
*      ; ADDITION MULTIPLY.
*      STF R2,++R1(R1)
*      ; SAVE 1ST ADDITION MULTIPLY IN UPPER
*      ; 1/2 OF BUTTERFLY.
*
*      END_2ND_LOOP:
*
*      STF R3,++R3(R1) ; SAVE 2ND ADDITION MULTIPLY IN UPPER
*      ; 1/2 OF UPPER BUTTERFLY.
*
*      END OF FINAL BUTTERFLY STAGE LOOP.
*
*      BIT REVERSE ADDITION LOOP SERIES.
*
*      THIS LOOP SERIES DOES ALL OF THE BIT REVERSE ADDITIONS AT THE END OF FAST
*      COLUMN TRANSFORM.
*
*      LDI 2,I,RO
*      LDI #46,ARI
*      ADDI 4,ARI
*      LDI #41,AR2
*      LDI 8,I,RI
*
*      LAST_OUTSIDE_LOOP:
*
*      LDI #R2,AR4 ; UPDATE POINTERS AND COUNTERS.
*      LSH -1,AR5

```

```

*
*      AR5,RC ; SET UP REPEAT COUNTER.
*      #R2++(I,RO),B,++R4++(I,RO),RO ; DATA POINTER UPDATE.
*      ARI,RI ; USE INITIAL ARI VALUE AS INNER LOOP
*      ; CONTROL.
*
*      SUBI 1,RC
*      #R4++(I,RO),B
*      LDI #R2,AR3 ; CONTINUE UPDATING POINTERS.
*
*      RPTB END_INSIDE ; TWO ADDITIONS ARE DONE IN EACH LOOP.
*
*      LAST_INSIDE_LOOP:
*
*      ADDF3 #R1,++R2++(I,RI),R0 ; ADD FIRST TWO DATA.
*      ADDF3 #R2,++R4++(I,RI),R1 ; ADD SECOND TWO DATA.
*      STF R0,++R1++(I,RI) ; SAVE FIRST ADDITION.
*
*      END_INSIDE:
*
*      STF R1,++R3++(I,RI) ; SAVE SECOND ADDITION.
*
*      END OF INSIDE LOOP FOR LAST LOOP SERIES.
*
*      ADDF3 #R1++(I,RO),B,++R2++(I,RO),B,RO ; UPDATE DATA POINTERS.
*      ADDF3 #R3++(I,RO),B,++R4++(I,RO),B,RO
*      ADDF3 #R3++(I,RO),B,++R4++(I,RO),B,RO
*      ADDF3 #R1++(I,RO),B,++R2++(I,RO),B,RO
*      CPTI R4,AR4 ; IS THIS LOOP COMPLETE?
*      BNEI LAST_INSIDE_LOOP ; DELAYED BRANCH, IF NOT.
*      LDI #R5,RC ; SET UP REPEAT COUNTER.
*      SUBI 1,RC
*      OR 0100H,ST ; SET REPEAT MODE.
*
*      BRANCH DELAYED TO LAST_INSIDE_LOOP.
*
*      RPTB LAST_BLOCK ; SINCE THERE ARE AN ODD NUMBER OF
*      ADDF3 #R1,++R2++(I,RI),R0 ; ADDITIONS, THE FINAL ONES ARE
*      ; DONE NOW.
*
*      LAST_BLOCK:
*
*      STF R0,++R3++(I,RI) ; SAVE ADDITION.
*
*      END OF LAST REPEAT BLOCK.
*
*      LSH 1,I,RO ; MULTIPLY IRO BY 2.
*      ADDI I,RO,RI ; UPDATE INNER LOOP CONTROL REGISTER.
*      CPTI 1,AR5 ; ARE CALCULATIONS COMPLETE ?
*      BGTI LAST_OUTSIDE_LOOP ; DELAYED BRANCH, IF NOT.
*      LDI #R4,AR2 ; UPDATE DATA POINTERS.
*      LDI R4,ARI
*      LSH 1,I,RI ; MULTIPLY IRI BY 2.
*
*      DELAYED BRANCH TO LAST_OUTSIDE_LOOP.

```

```

* END OF LAST LOOP SERIES.
*
* MULTIPLY COEFFICIENT ZERO BY .5, IF NOT ZERO.
*
*      LDF  #R6,R0      ; SET ZERO FLAG IF #R6 = 0.
*      BEQD DONT_STORE ; IF COEFFICIENT IS ZERO, DON'T DO
*           ; THIS.
*      LSH  24,#R5     ; USE INTEGER PART FOR FLOAT DIVIDE
*           ; BY 2.
*      SUB13 #R5,#R6,#R1
*      NOP
*
* DELAYED BRANCH FROM HERE IF VALUE IS NOT TO BE STORED.
*
*      STI  #R1,#R6    ; STORE, IF EXPONENT MASK'N'T -128.
*
* DONT_STORE:
*
*      RETS

```







## Appendix E3. FCT Cosine Tables File

```
*
* APPENDIX E3
*
* FCT COSINE TABLES FILE
*
* TO BE LINKED WITH FCT SOURCE CODE FOR 32 POINT FCT.
*
* COEFFICIENTS ARE  $1/(2 * \cos(N*PI/2M))$ , WHERE N IS A NUMBER FROM 1 to
* M-1. M IS THE ORDER OF THE TRANSFORM.
*
* FOR A 32 POINT FCT, N IS IN THE FOLLOWING ORDER:
*     1, 15, 3, 13, 5, 11, 7, 9,
*     2, 14, 6, 10,
*     4, 12,
*     8
*
* THE LAST VALUE IN THE TABLE IS 2/M.
*
*
*     .global  COS_TAB
*     .global  M
*
M     .set     16
*
*     .data
*
COS_TAB
*     .float   0.5024193
*     .float   5.1011487
*     .float   0.5224986
*     .float   1.7224471
*     .float   0.5669440
*     .float   1.0606777
*     .float   0.6468218
*     .float   0.7881546
*     .float   0.5097956
*     .float   2.5629154
*     .float   0.6013449
*     .float   0.8999762
*     .float   0.5411961
*     .float   1.3065630
*     .float   0.7071068
*     .float   0.1250000
*     .end
```

## Appendix E4. Data File

```
*
* APPENDIX E4
*
* DATA FILE
*
*       .global COEFF
*
*       .data
*
COEFF
    .float 137.0
    .float 249.0
    .float 105.0
    .float 217.0
    .float 73.0
    .float 185.0
    .float 41.0
    .float 153.0
    .float 9.0
    .float 121.0
    .float 233.0
    .float 89.0
    .float 201.0
    .float 57.0
    .float 169.0
    .float 25.0
    .end
```

## **Appendix F. Test Vectors, 64-Point Sine Table, Link Command File**



# Appendix F1. Example of a 64-Point Vector to Test the FFT Routines

\*  
\* APPENDIX F1  
\*  
\* EXAMPLE OF A 64-POINT VECTOR TO TEST THE FFT ROUTINES  
\*

X =

0.2113  
0.0624  
0.7599  
0.0087  
0.8096  
0.8474  
0.4324  
0.8075  
0.4832  
0.6135  
0.2749  
0.6807  
0.6538  
0.4899  
0.2741  
0.7626  
0.9933  
0.8360  
0.7469  
0.0378  
0.4237  
0.2613  
0.2403  
0.9405  
0.1167  
0.6290  
0.3510  
0.9550  
0.4943  
0.0365  
0.2260  
0.8159  
0.2284  
0.6553  
0.6621  
0.7075  
0.2408  
0.6907  
0.1062  
0.2640  
0.7034  
0.4021  
0.6553  
0.9700  
0.0380  
0.0988  
0.2560

0.5598  
0.7166  
0.1402  
0.7054  
0.0178  
0.2611  
0.1358  
0.0503  
0.5782  
0.2432  
0.8448  
0.3876  
0.7256  
0.2849  
0.6767  
0.8642  
0.1943

\* 64-POINT FFT CORRESPONDING TO VECTOR X

Y =

30.3774  
1.7780 - 2.5364i  
-1.0376 - 2.9799i  
-1.0123 + 2.4889i  
0.6594 + 2.3639i  
-1.5228 - 0.7827i  
-3.8171 - 0.2050i  
-2.7096 + 1.2841i  
2.1822 - 1.6863i  
0.2879 + 1.8671i  
-1.5479 + 1.6298i  
-0.6366 - 0.1176i  
2.2902 + 1.5249i  
-2.4837 - 0.5842i  
-1.7338 + 0.0738i  
-0.2180 - 0.4726i  
-0.2104 + 0.4897i  
-1.7473 - 1.0213i  
0.1233 - 2.3915i  
-0.6415 - 1.1144i  
-2.7719 - 0.4802i  
-0.0063 - 0.3885i  
-0.7163 + 1.5662i  
0.3218 - 1.3316i  
-0.7823 + 1.0607i  
-0.2533 + 2.8270i  
-1.0813 - 2.7861i  
3.4619 + 1.9485i  
3.0382 + 1.3853i  
3.2699 + 2.3564i  
-1.9511 - 0.7714i  
1.8753 + 0.2867i

-1.5474  
1.8735 - 0.2867i  
-1.9511 + 0.7714i  
3.2099 - 2.3564i  
3.0532 - 1.3695i  
3.4869 - 1.9465i  
-1.0813 + 2.7861i  
-0.2953 - 2.8270i  
-0.7823 + 1.0607i  
0.3218 + 1.3316i  
-0.7163 - 1.5682i  
-0.0063 + 0.3885i  
-2.7719 + 0.4802i  
-0.6415 + 1.1144i  
0.1233 + 2.3915i  
-1.7473 + 1.0213i  
-0.2104 - 0.4897i  
-0.2180 + 0.4726i  
-1.7338 - 0.0738i  
-2.4837 + 0.5842i  
2.2902 - 1.5549i  
-0.6366 + 0.1176i  
-1.5479 - 1.6298i  
0.2879 - 1.8671i  
2.1622 + 1.6833i  
-2.7096 - 1.2841i  
-3.8171 + 0.2658i  
-1.5228 + 0.7327i  
0.6594 - 2.3359i  
-1.0123 - 2.4889i  
-1.0376 + 2.3999i  
1.7780 + 2.5584i

# Appendix F2. File to Be Linked with the Source Code for a 64-Point, Radix-4 FFT.

```

*
* APPENDIX F2
*
* FILE TO BE LINKED WITH THE SOURCE CODE FOR A 64-POINT, RADIX-4 FFT.
*
*
*      .globl  SINE
*      .globl  N
*      .globl  H
*
*      .set   64
*      .set   6
*
*      .data
*
*      .float 0.000000
*      .float 0.098017
*      .float 0.195090
*      .float 0.290285
*      .float 0.362483
*      .float 0.471397
*      .float 0.535570
*      .float 0.634393
*      .float 0.707107
*      .float 0.773010
*      .float 0.831470
*      .float 0.881921
*      .float 0.923880
*      .float 0.956940
*      .float 0.980785
*      .float -0.995185
*      .float -1.000000
*      .float -0.995185
*      .float -0.980785
*      .float -0.956940
*      .float -0.923880
*      .float -0.881921
*      .float -0.831470
*      .float -0.773010
*      .float -0.707107
*      .float -0.634393
*      .float -0.535570
*      .float -0.471397
*      .float -0.362483
*      .float -0.290285
*      .float -0.195090
*      .float -0.098017
*      .float 0.000000
*      .float 0.098017
*      .float 0.195090
*      .float 0.290285
*      .float 0.362483
*      .float 0.471397
*      .float 0.535570
*      .float 0.634393
*      .float 0.707107
*      .float 0.773010
*      .float 0.831470
*      .float 0.881921
*      .float 0.923880
*      .float 0.956940
*      .float 0.980785
*      .float 0.995185
*
*
*      .float 1.000000
*      .float 0.995185
*      .float 0.980785
*      .float 0.956940
*      .float 0.923880
*      .float 0.881921
*      .float 0.831470
*      .float 0.773010
*      .float 0.707107
*      .float 0.634393
*      .float 0.535570
*      .float 0.471397
*      .float 0.362483
*      .float 0.290285
*      .float 0.195090
*      .float 0.098017
*      .float 0.000000
*      .float -0.098017
*      .float -0.195090
*      .float -0.290285
*      .float -0.362483
*      .float -0.471397
*      .float -0.535570
*      .float -0.634393
*      .float -0.707107
*      .float -0.773010
*      .float -0.831470
*      .float -0.881921
*      .float -0.923880
*      .float -0.956940
*      .float -0.980785
*      .float -0.995185
*
*
* COSINE

```

## Appendix F3. Link Command File

```
*
* APPENDIX F3
*
* LINK COMMAND FILE
*
* DO NOT TYPE IN THESE FIRST SEVEN LINES
-o 12opt64.out
12fopt.obj
sin64.obj

SECTIONS
{
    .text : {}
    .data : {}
    IN 809800h : { 12fopt.obj(IN) }
    .bss 809C00h: {}
}
```