



**AP-618**

**APPLICATION  
NOTE**

# **Software Concerns of Implementing a Resident Flash Disk**

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# SOFTWARE CONCERNS OF IMPLEMENTING A RESIDENT FLASH DISK

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## PURPOSE

The purpose of this Ap-Note is to detail the software aspects of implementing an Intel Resident Flash Array (RFA) as a Resident Flash Disk (RFD) in an embedded system with commercially available Flash Translation Layer (FTL) software.

An RFD can help mitigate many issues that constrain the storage subsystem for embedded systems. Those issues include:

- Embedded systems must operate in harsh environments; they are dropped, banged, vibrated, overheated, etc.:
  - A Flash RFD is rugged and non-volatile because it is solid-state and has no moving parts.
- Embedded systems must have high performance:
  - A Flash RFD, because it is basically memory, not rotating magnetic media, has extremely fast access capability.
- Embedded systems can be battery operated and must be low power:
  - A Flash RFD is very low power as it is silicon and has no motors to spin, or servos to move. The standby/deep dower-down mode current draw of a flash part is typically measured in micro-amps ( $\mu$ A).
- Embedded systems may be limited in size:
  - A Flash RFD requires very little board space and is tiny compared to even the smallest hard drive.
- Embedded systems tend to be price sensitive:
  - A Flash RFD is relatively inexpensive.
- Embedded systems have to be very flexible:
  - A Flash RFD is adaptable to just about ANY situation.

For the purposes of this document, we will use a “real-life” example based on the Intel “Ultra Low Power” (ULP) Intel486™ SX Processor Evaluation Board and M-Systems TrueFFS\* FTL flash software for DOS.

Familiarity with the PCMCIA PC Card Standard is helpful for understanding the software pieces of our example. Please see the Reference section for information on contacting PCMCIA in order to obtain a copy of “The Standard.”

## INTRODUCTION

The exacting requirements for today’s embedded processor platform market makes extreme demands upon all of the associated embedded sub-systems. Among those are the data and code storage functions. Intel’s Flash components are well suited to implement a ruggedized, low-power, low-cost storage system, in the form of a Resident Flash Disk (RFD).

A RFD is basically a Resident Flash Array (RFA), which is one or more flash components typically designed onto the base system, used with software to make it operate as a drive.

The flash that comprises the RFD has the unique capability of non-volatility (like ROM) combined with in-place erase and write (similar to RAM). With the proper software, Intel Flash can function as a solid-state disk drive, without exhibiting the disadvantages inherent in a rotating/magnetic-media drive. In general, as well as for the purposes of this Ap-Note, the proper software is FTL. Although there are quite a few sources for FTL, we base our example on M-System’s<sup>1</sup> TrueFFS\* FTL<sup>2</sup>. Please keep in mind that though the details of the implementation will be specific to TrueFFS, the general concepts and functions will apply to virtually any of the available FTLs or other flash managers and/or flash filing systems. This also applies for the example target platform. For our example, the specific target platform is the Intel Ultra Low Power (ULP) 486 SX Evaluation Board. However, the general concepts presented here will apply to virtually any embedded system.

<sup>1</sup>M-Systems provides a full range of Turn-key FlashDisk solutions that uses TrueFFS, are plug and play and uses Intel advanced Flash components. These solutions include ISA Bus and PC104 FlashDisks, and the DiskOnChip\* which is a Plug & Play FlashDisk module in a standard 28/32 pin DIP JEDEC compatible package.

<sup>2</sup>Some of the information in this document is taken from M-System’s TrueFFS documentation.

## SYSTEM REQUIREMENTS

We are using the ULP 486 SX for our example because it is a high performance demo platform which exhibits characteristics that are pertinent to a wide range of embedded applications. Its compact size, high-degree of functionality, high performance, extreme flexibility, availability of many development environments and op-

erating systems, and PC desktop/notebook compatibility make it easily adapted to the many and varied requirements placed on embedded systems. It also contains a considerable quantity and variety of flash, including a simple RFD which is well suited for use with M-Systems TrueFFS FTL. Figure 1 shows the general layout of the ULP 486 SX

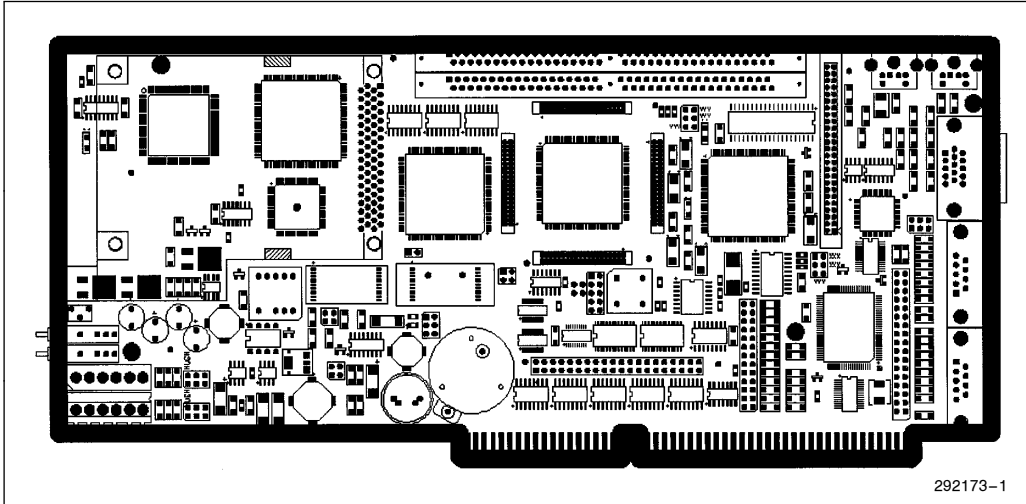


Figure 1. Intel486™ SX ULP Evaluation Board

The following chart gives an operational overview of our target embedded platform. Included is an assessment of how our specific features might apply to a different implementation.

Item	This Specific Example	Your Application
Processor:	Intel486™ SX ULP (80x86)	Any processor can be used, however a version of FTL for your particular processor will be required.
Flash:	RFD- 1 x Intel 28F008	Varies. Could be an RFD based on any variety or quantity of Intel Flash components, or it could be a Flash Card plugged into a PC Card socket.
Flash Media Manager:	FTL (M-System's TrueFFS*)	Any FTL. The low level routines of the FTL you select will probably be different but will have similar requirements and functions.
Operating system:	DOS-FAT	A "sector" based filing system of some sort is required for most FTLs. The FTL typically needs to be "ported" to your specific operating system (if other than DOS).
Other Flash:	PC Card Socket (Cirrus 6710)	Varies. A typical system will usually only implement a single "variety" of flash. A PC Card socket accepts flash in a Flash Card.
	XD Simm (not covered)	Varies. If your application would benefit from execute in place (XIP) on the memory bus, then you might consider using this style of flash. Please contact your Intel representative for more information on XIP and Intel XD Flash.
Low Level FTL Interface:	Socket Services Like Interface	Varies depending on the particular FTL you decide to implement in your design. However, most of the elements of, and all of the functionality described here in the Socket Services like routines, are typically required in the low level routines of any flash undertaking.





## WHY DO I NEED SOFTWARE FOR FLASH?

One of the unique characteristics of flash technologies is the typical requirement that a flash cell be erased before it can be written.<sup>3</sup> For Intel Flash, generally,<sup>4</sup> a flash cell starts as erased to a 1 (one), and can be written/programmed to become a 0 (zero). Early flash chips were organized such that the entire chip was erased at one time (called “bulk erase”). Newer flash chips are organized into erase blocks—typically 64 KBytes<sup>5</sup> in size. Arranging the flash chip into erase blocks provides a mechanism for software to manage the data stored on the chip. Keep this in mind, as there will be a quiz later.

Traditional data storage programs, such as DOS and other file systems, are designed to update or rewrite data in place. Here is a specific example: think of how DOS deals with the File Allocation Table (FAT) when you write a file. It takes a specific sector on the media and rewrites it with the new or changed data concerning the file. Add or change a file and the sector gets changed. The same sector or sectors are constantly rewritten with new or different FAT data. This is a reasonable mode of operation for rotating magnetic media but at odds with the basic operation of flash. Another difference between drives and flash is the traditional drive sector (512 bytes is typical) tends to be quite a number of bits smaller than the typical flash block.

Flash management software can solve these issues in a number of ways, two of which follow:

1. When rewriting any sector (512 bytes of data typical), the flash manager can perform the following steps:
  - a. Save off the rest of the “good data” stored in the larger (64 KB typical) flash erase unit. This will be comprised of the other valid “sectors” in the flash block. It will have to put the good data “somewhere” else such as a spare flash block.

- b. Next, erase the current block where the data wants to be rewritten.
- c. Write in the new data.
- d. Restore the rest of the good data from the temporary storage.
- e. Finally clear out the temporary storage location (i.e. erase the spare flash block).

Whew, this sure looks and sounds tedious and highly inefficient! No matter what the style of flash, a purely erase before write type flash management scheme is definitely NOT the best overall performance solution.

2. When rewriting a sector it can:
  - a. Write the updated data to another free/erased portion of the flash.
  - b. Then point a lookup entry, or translation table to the other location in flash
  - c. Finally, the old portion of flash is marked as “dirty” or deleted for later clean-up.

The second method is what, in general, the FTL software does for flash. This is the most write efficient, highest performance and most highly recommended method for handling flash.

## FTL IS IT!

Flash Translation Layer is a robust, widely accepted, industry standard flash manager. It is widely available from a multitude of sources. Its function is to take disk drive specific software requests and convert/translate them to flash media accesses. It handles flash blocks by creating small, virtual, sector-sized (usually 512 byte) blocks out of the larger flash blocks. Additionally, FTL handles the “special needs” of flash such as handling the special read, write, and erase requirements that flash technology exhibits. Also, many, if not all, of the varieties of FTL offer a boot solution. This makes the operating system loadable from flash as if it were

<sup>3</sup>See the section on Bit Twiddling for the exception to this “rule”.

<sup>4</sup>Current flash technology (for example i28F008 and i28F016) uses 1 flash cell per bit. Intel has announced a technology that in the near future will be able to represent more bits per flash cell using a technology known as Multi-Level Cell or MLC. The concepts presented here may not all apply to MLC technology.

<sup>5</sup>Flash block size and organization may vary from component to component- please refer to the data-sheet for the particular flash component you are using. In general, RFA/RFD implementations use Intel FlashFile™ components (i28F008, i28F016, etc.) which are organized in symmetrical 64KB and 128KB erase blocks. Boot Block components can be used but pose extra burdens on the erase and write routines in order to adapt to the non-symmetrical erase units present in the components.



loading from a more common bootable device such as a floppy or hard disk (more on this in a later section).

In the next section, we will describe generally how FTL works.

### FTL FUNCTIONAL OVERVIEW: WHAT IS A FLASH TRANSLATION LAYER?

As we have said, FTL is a sector based flash manager that provides logical to physical sector translation. Thus the name Flash Translation Layer. FTL performs sector mapping to allow Flash to appear as a drive-like, sectored, rewrite “in place” type storage media. While the host file system sees the Flash card or resident Flash array (RFA) as a continuous sectored medium, FTL relocates these sectors transparently to the operating system, and tracks the logical-to-physical relationship. Figure 2 provides a simplified graphical representation of the sector translation and relocation that occurs with FTL. The MAP reference in the diagram is explained in greater details in the “FTL Functional Details” section of this document.

This logical-to-physical mapping allows the Operating System to concern itself with only file operations. Because the O/S already oversees these file operations, the FTL solution can provide compatibility with existing

applications and media utilities while presenting a small code footprint. Of important note here is the fact that FTL is able to relocate sectors to any position in the Flash media, making large Flash blocks appear as smaller erasable sectors. This is basically what FTL does to earn its keep.

### WHERE DOES FTL FIT IN THE SOFTWARE SCHEME?

FTL, for the most part, needs to intercept the data store/retrieve requests somewhere between the software wanting to talk to the drive and the low level routines that specifically read or write data to the drives. In a x86 PC like system, and on other types of platforms, there are typically two logical places to grab the drive access routines in order to handle the drive requests actually being serviced by flash. Each has its advantages and disadvantages which we will discuss.

### Block Device Driver

One of the most common ways to intercept the normal drive handlers is to have the operating system load a Block Device Driver which inserts itself in the lower levels of the operating system. M-System’s FTL TFFS.COM and TFFSCS.COM are two examples of a

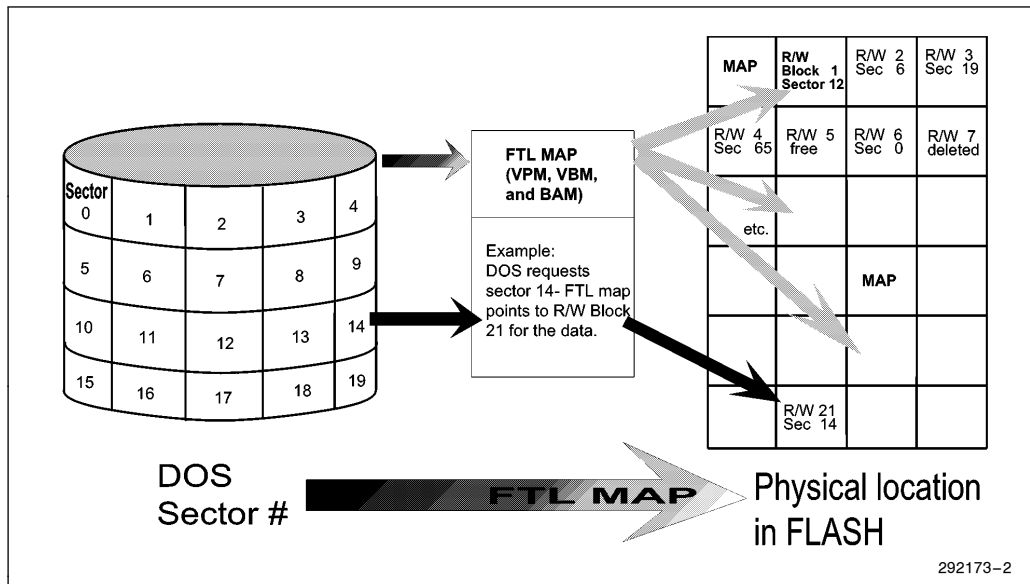


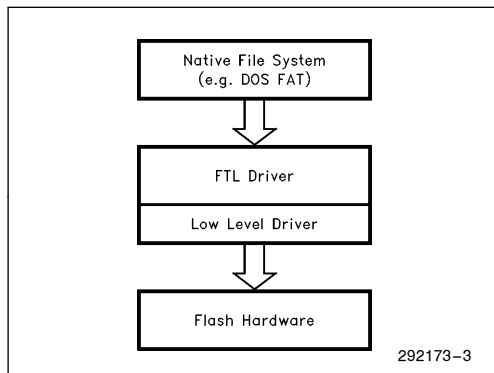
Figure 2. FTL Sector Relocation

FTL Block Device Driver. The major advantage to using the Block Device Driver approach is that it is specific to the operating system and therefore can have better knowledge about how the filing system works. Because it knows how the operating system works, it can use performance enhancement tricks specific to that operating system. One such trick used by M-System's TrueFFS DOS FTL driver is a technique called "FAT Snooping." This allows the FTL to monitor writes to the File Allocation Table to know when DOS has deleted a sector thus allowing the FTL to perform clean-up more efficiently.

The major disadvantages to grabbing the drive control at this level are:

- The driver must be operating system specific. This requires that a new driver be developed for each operating system. (Hmmm, this is also an advantage- let's not go into an explanation on the inherent duality of nature!)
- You cannot boot from the flash as, typically, the device drivers, including the FTL block device drivers, are loaded AFTER the operating system starts. You can't read from the flash to boot as the driver has to be loaded by the operating system you are trying to load. We will go into more boot issues in the next section, as well as later in this paper.
- Software that bypasses the operating system to talk to the low level Int13h BIOS routines OR talks directly to the drive controller hardware will not work with a flash/FTL type drive.

Figure 3 gives a general layout of this layered approach, where FTL works with the existing File System to control the Flash media.



**Figure 3. FTL Block Device Driver Software Layers**

### BIOS Interception Driver

The other location to intercept drive requests in order to have flash handle the data is in the lowest levels of the system itself. In the vast majority of systems, that lowest level software is something known as BIOS. BIOS stands for Basic Input/Output System. It provides the services to handle the lowest level interface tasks to the various standard I/O systems like drives, serial ports, parallel ports, keyboard, etc. In x86 PC-like systems, the INT13h handler is the one responsible for servicing the system drives (HDD and FD) and it is this software vector that needs to be intercepted by an FTL. M-Systems provides a driver called TFFSBIOS.EXB which is loaded as an Expansion ROM and intercepts INT13h.

The advantages to this approach are:

- Boot capable—the BIOS uses INT13h BIOS calls to boot so if the FTL/flash device is set up as one of the standard boot drives then the OS will look to the FTL partition on the flash to try and boot.
- Software that bypasses the operating system and interfaces directly to the low level BIOS INT13h routines will still work on flash.

Disadvantages:

- Flash Software must be present in non-volatile memory such as system BIOS flash. This is more complicated to deal with than a driver stored on a disk drive and loaded to memory.
- No performance tricks such as snooping can be performed as we do not know what will be accessing us.
- The software that accesses the drive must be "well behaved" in order for the interception to work. Software that tries to talk directly to a HDD controller will NOT work even with a BIOS Intercept driver.

Now that we have seen the Why and Where FTL works, let's take a look at the specific details of HOW it performs its magic.



## FTL FUNCTIONAL DETAILS—THE GUTS OF THE BEAST

In these sections, in reasonable technical depth, we will describe the details of the FTL structures as defined by the industry standard FTL specification (available from PCMCIA).

### FTL Format Overview

When a FTL format occurs, all of the FTL overhead, such as the Virtual Block Maps (VBMs), FAT table, Root Dir Entries, and any other structures, is put in the

first two logical blocks. Figure 4 illustrates how multiple flash chips in a flash array might be organized with the FTL structures after a format.

When writing a new sector or changing an existing one, FTL looks for available (erased) flash in the form of an unallocated Read/Write Block, writes the data, and sets up the pointers and maps accordingly. Sectors containing valid data for a specific file can be scattered anywhere on the flash memory in this fashion. You never know exactly where the data will go physically in the flash, but rest assured that you will definitely get it back when you ask for it.

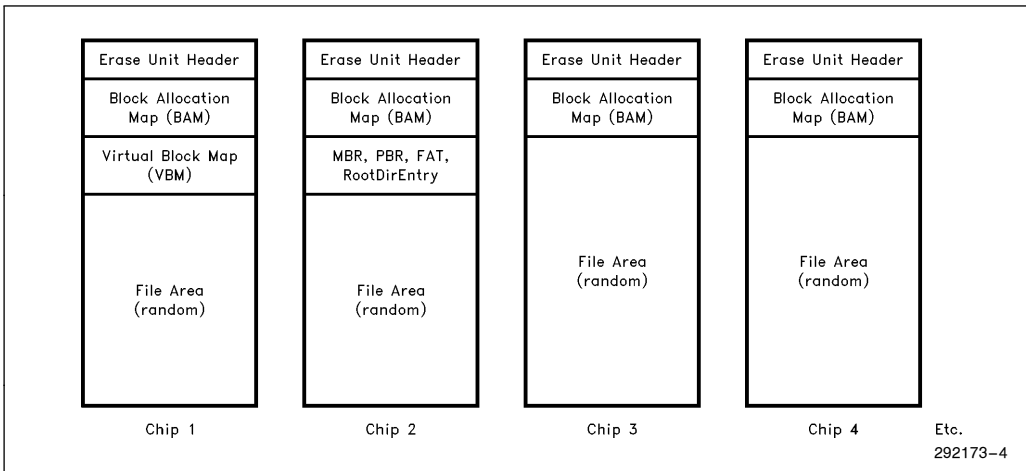


Figure 4. Typical FTL Overhead Organization

### Erase Unit Header and Block Allocation Information

For allocation purposes, each Erase Unit is evenly divided into arrays of Read/Write Blocks of equal size. For the purposes of DOS, and also to appear more

“drive-like”, the Read/Write blocks tend to be 512 bytes in length. You can also think of them as sector sized.

Figure 5 illustrates the multiple Read/Write Block per Erase Unit concept.

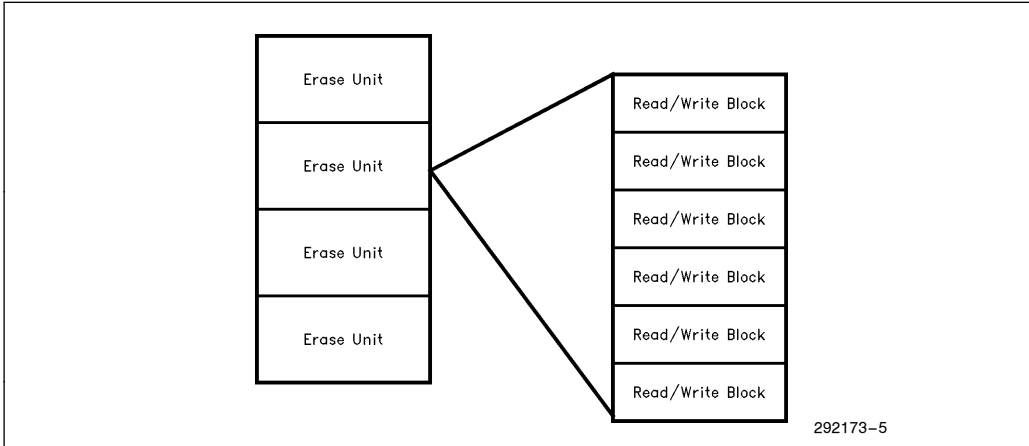


Figure 5. Erase Unit Organization



The size of a Read/Write block is the same as a virtual block as viewed by FAT. A Read/Write block used to store a FTL overhead structure is called a Control Unit. The BAM, VBM, EUH, etc. all reside in various Control Units. The first Read/Write block of each Erase Unit is a Control Unit which contains the Erase Unit Header (EUH) for that Erase Unit. It includes specific information about the Erase Unit and global information about the format of the FTL partition. Table 1 details the contents of a typical EUH.

**Table 1. Erase Unit Header Fields**

Offset	Field
0	LinkTarget Tuple
5	DataOrganization Tuple
15	Number of Transfer Units
16	Erase Count
20	Logical Erase Unit (LUN) Number
22	Read/Write (sector) Size
23	Erase Unit Size(in log2 form)
24	First Physical Erase Unit of Start of Partition
26	Number of Erase Units
28	Formatted Size
32	First Virtual Map Address on the Media
36	Number of Virtual Map Pages
38	Flags
39	Code
40	Serial Number
44	Alternate Erase Unit Header Offset
48	BAM Offset
52	Reserved

Each Erase Unit also contains allocation information for all of the Read/Write Blocks within the unit.

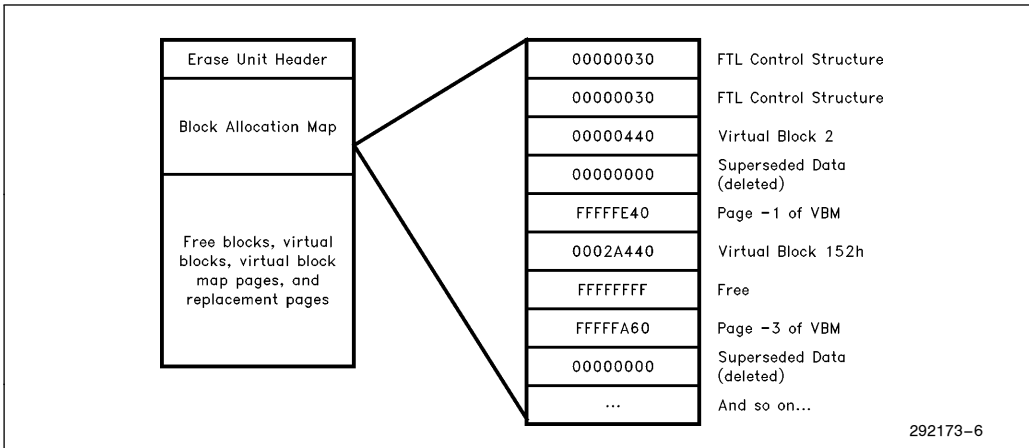
For each Read/Write Block, a 4-byte value tracks the block's current state. This is the case for all Read/Write blocks in that Erase Unit. This section, comprised of one or more Control Units, is located right after the EUH Control unit and is called the Block Allocation Map or BAM. At any point in time, a Read/Write Block in an Erase Unit may be free, deleted, bad or allocated. Table 2 details the exact status values used to indicate those states.

**Table 2. BAM Status Values**

Value	Meaning
FFFFFFFF	Free
00000000	Deleted
00000070	Bad
00000030	Control
xxxxxx40	Data of Map Page
xxxxxx60	Replacement Page

In Figure 6, we show an example of what a BAM might look like after the FTL partition has been used.





**Figure 6. BAM Example**

The BAM entries for Virtual Block Maps are negative numbered in contrast to the BAMs for Virtual Block data which are positive numbered. This is the only way to distinguish between the two. In the example in Figure 6, the 00000440 in the BAM is the virtual block data number 2 (each block is 200 hex bytes) while the FFFFFFFE40 BAM entry is the last page of the VBM.

The number of Control Units used depends on the size of the BAM which depends on the ratio of Erase Unit size to the Read/Write block size

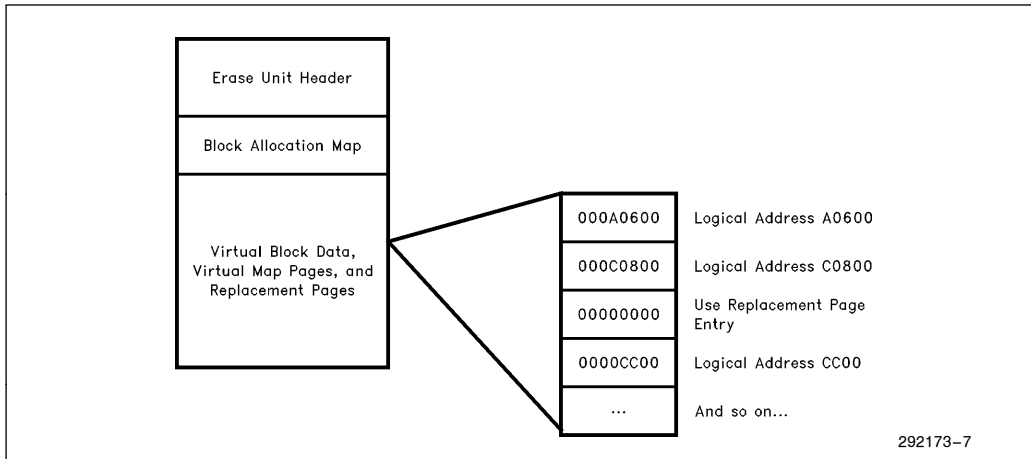
**Virtual Block Map**

The FTL uses a data structure known as the Virtual Block Map (VBM) to map requests for sectors from the

higher level software layers to the flash. The file system *thinks* it is requesting physical sectors but through FTL, they are actually virtual sectors, also known as Virtual Blocks. FTL maps the Virtual Block requests through the VBM to obtain a logical address. The physical address on the flash is then determined from the logical address and the data is returned to the filing system.

The VBM is comprised of an array of 32-bit entries, each of which represents a logical address on the media where a Virtual Block’s data is stored. The Virtual Block number requested by the higher level software layer is used as an index into this array. Please see Figure 7 for an example of a VBM.





**Figure 7. Virtual Block Map Example**

The VBM is subdivided into pages. Each page of the VBM is the same size as a Virtual sector of the FAT file system. Since each entry in the VBM is 4 bytes, each page holds (Virtual Blocks/4) number of entries, and from that we can figure out how much virtual space each map represents and how many pages we need to map the whole virtual space.

Space is always reserved on the media to store a VBM large enough to track the allocation of all the Virtual Blocks on the flash. However, when the flash is formatted, the FTL may choose to keep only a portion of the VBM on the media, and store the rest of it in RAM. The amount of VBM stored on the media is indicated by the FirstVMAddress field of the Erase Unit Header (see Table 1). If the first VMAddress is set to 0, the FTL maintains all of the VBM entries on the media. If the FirstVMAddress exceeds the FormattedSize, none of the VBM entries are maintained on the media by the FTL.

When all or a portion of the VBM is not maintained on the media, it has to be reconstructed in RAM every time the system is booted up running FTL, or, as in the case of a PC Card slot with flash cards, every time a

card is re-inserted. FTL uses the BAM information to fill out the entries of the VBM in RAM.

If a VBM entry is all ones, the Virtual Block does not exist on the media. If it is all zeroes, the logical address of the Virtual Block is described on a Replacement Page.

### Replacement Pages

Each page of the Virtual Block Map may have a Replacement Page. Values in a Replacement Page override entries in the original VBM pages. Replacement pages are allocated from free Read/Write Blocks in any Erase Unit. The FTL locates allocated Replacement Pages by scanning the block allocation information on the media. This scan may be performed when the media is inserted in the host system or when a VBM entry of zero is encountered. Replacement Pages cannot themselves be replaced. The block allocation information entry for a Replacement page uses the same virtual address as the original VBM page. FTL distinguishes between the two by looking at the last byte of the entry: 40H for VBM pages and 60H for Replacement Pages.



### SOCKET SERVICES—WHAT IS IT AND WHY DO I NEED TO USE IT?

The lowest level of our example RFD software solution is a PC Card-like Socket Services (SS) subset. The illustration in Figure 8 gives a general overview of the details of the full PC Card standard software layers.

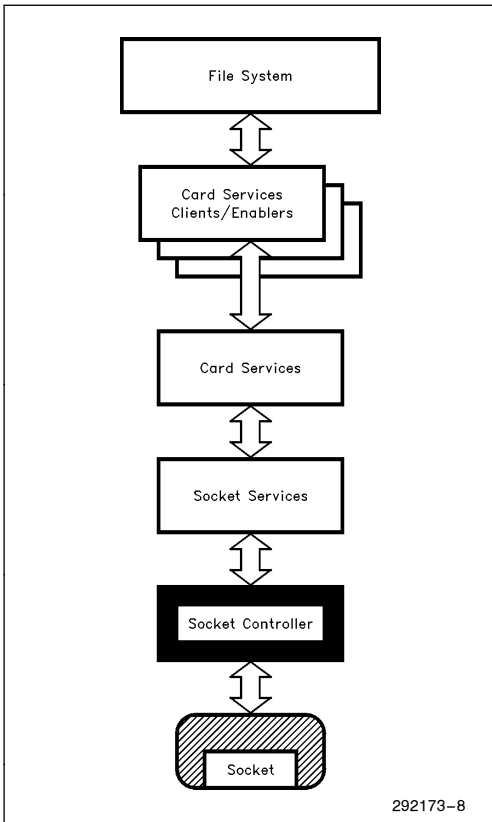


Figure 8. PC Card Software Suite General Layout

The Socket Services (SS) standard gives an industry recognized, standardized, interface method for other software (i.e. the FTL core) to interact with the low level routines. This separation of the low level routines from the other PC Card software allows the often different and hardware specific portion of the code to be cleanly separated from the (mostly) unchanging higher level functions. Please refer to Appendix B for a listing of the full PC Card Socket Services functions for reference.

Other FTL implementations may organize the low level interface differently than the SS interface we describe here but will typically have similar requirements, and utilize similar sorts of functions. The following figure 9 illustrates the software layers for our specific implementation.

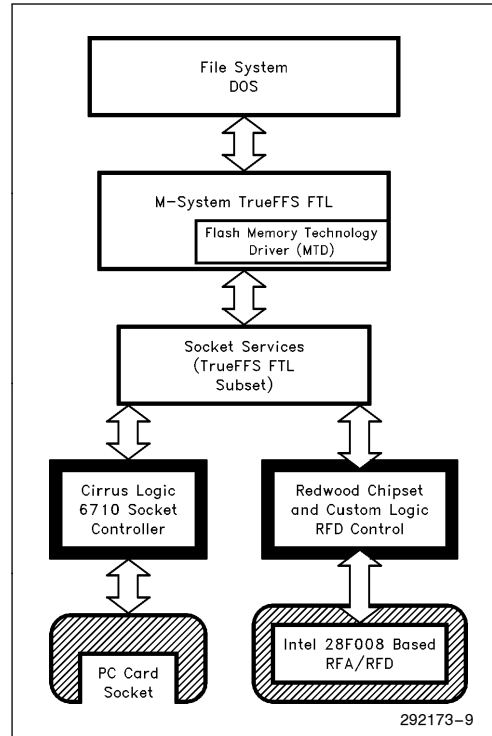


Figure 9. M-System's FTL Software Layers

### Implementation Issue: Real-Mode Flash Memory Sliding Window

There is a technical requirement of the example FTL that is difficult to describe but relatively easy to implement. It is something called a real-mode sliding window to the RFD. "A real-mode sliding what?" you might be asking yourself right now. Basically, this a method for accessing the flash memory of the RFD that is mapped "somehow" into the lower one megabyte region of memory. That region is known as the real-mode, or DOS memory region. Typically, some form of hardware such as the main system chipset, special logic (like a PLD), or a combination of the two, will cause the RFD to be accessible in a real-mode part

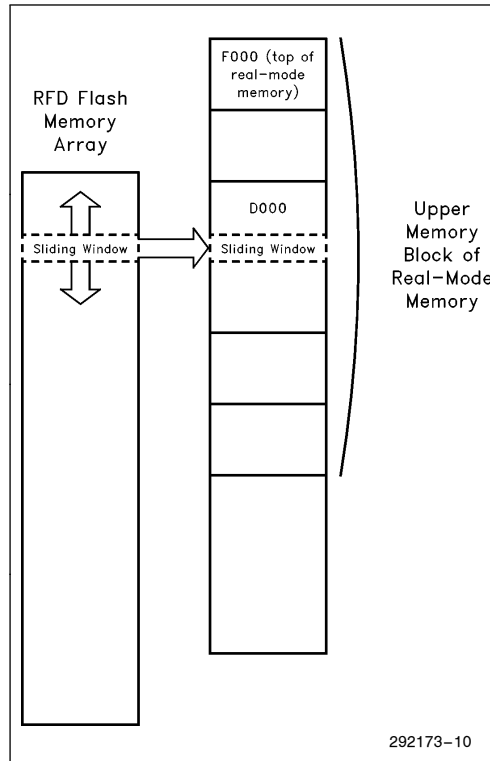


of memory. That part of memory is usually in the upper memory block (UMB) region (A000:0 FFFF:F) starting in the D000 page. For the ULP platform, a special capability of the Cirrus Logic PicoPower Redwood chipset is used along with logic in a Xilinx PLD to do just that. "But wait a minute, what about the sliding part?" Ahhh, here is the tricky part. The flash memory size of the RFD is usually quite large in order to implement a reasonable size flash drive. The range of RFD sizes will normally start at a minimum one megabyte, single 28F008 8 megabit FlashFile component. The ULP RFD is this exactly. If we were to map the entire RFD into real-mode memory, the one megabyte of flash memory would take up the ENTIRE real-mode range leaving no room for DRAM, operating system (DOS), video RAM, system BIOS, other BIOS, etc! The trick is to only map into real-mode a small portion of the flash, or window, and then move the window around to be able to access the entire address range of the flash RFD. This concept is similar to the LIM/EMM memory scheme. Our example FTL is quite happy with a minimum 4 kilobyte window into the flash and will work well with that size or whatever the hardware is capable of supplying. The ULP logic controlled window can be programmed for 16K, 32K, or 64K windows but for the ULP implementation, we have settled on using 16K windows in order to conserve the UMB area. A window control register is used to pick where in the flash we are looking. The exact operation of the ULP registers will be discussed in detail in a following section.

**IMPORTANT:**

It is important to remember to exclude the real-mode UMB flash window from use by a memory manager. The window will not work if a memory manager, such as EMM386, has paged DRAM into the memory space needed by the sliding window!

The following Figure 10 attempts to illustrate the concept of the real-mode sliding window.



**Figure 10. Real-Mode Sliding Window**

The real-mode sliding window is the most common flash mapping technique as it eliminates the difficulties of accessing protected mode addresses from a DOS driver. In theory, accessing a RFD that is mapped in protected mode is possible, however, the majority of FTL solutions including the M-Systems FTL are limited to operating strictly from real-mode. If your micro-controller, embedded system and/or operating system enjoys the ability to address flash memory in a directly mapped, linear fashion, then the examples of windowing in the code listings that follow do not apply to you.

Now let's take a look at the requirements for the FTL socket services.

## Low Level: FTL Socket Services Subset

The Socket Services required for our example is a simple subset of the PC Card standard socket services. As you have seen before, Figure 3 shows the typical software layers of FTL and Figure 9 shows our reference software's specific layers. For the socket services/low level driver layer, there are only 11 functions out of the full PC Card standard 30 valid functions (refer to appendix B) that you need to support for the M-System's FTL. The required functions are as follows:

Name	Number	Description
<i>GetAdapterCount</i>	(80h)	Returns the number of adapters supported by all Socket Services handlers in the host system. Also used to determine if one or more Socket Services handlers are installed.
<i>GetSocket</i>	(8Dh)	Returns the current configuration of the socket identified by the input parameters.
<i>GetSSInfo</i>	(83h)	Returns the version/compliance level of the Socket Services interface supporting the adapter as specified by the input parameters. It also identifies the adapters serviced by the handler.
<i>GetStatus</i>	(8Fh)	Returns the current status of the card, socket, controls and indicators for the socket identified by the input parameters.
<i>GetWindow</i>	(88h)	Returns the current configuration of the window specified by the input parameters.
<i>InquireAdapter</i>	(84h)	Returns information about the capabilities of the adapter specified by the input parameters.

Name	Number	Description
<i>InquireSocket</i>	(8Ch)	Returns information about the capabilities of the socket specified by the input parameters.
<i>InquireWindow</i>	(87h)	Returns information about the capabilities of the window specified by the input parameters.
<i>SetPage</i>	(8Bh)	Configures the page specified by the input parameters. Only valid for memory windows. This service is unsupported by PC Card-32.
<i>SetSocket</i>	(8Eh)	Sets the current configuration of the socket identified by the input parameters.
<i>SetWindow</i>	(89h)	Sets the configuration of the window specified by the input parameters.

## FTL Socket Services: The CODE

In this section, we will give you the details of, and where to look for listings of the important parts of the actual code. This code, or a reasonable facsimile thereof, was used to implement the socket services for M-Systems FTL on the ULP Evaluation board. Please be aware that because the board also contains a Cirrus Logic 6710 Socket Controller, some of the code may be specific to it. The PC Card implementation of the ULP board does not fall into the scope of this Ap-Note, so we will be conveniently ignoring any such code here.

On the ULP board, the page register for sliding window control is located at I/O address 1096. The Window size and enable register is located at I/O address 2096. Watch for references to these ports as they will be important! Your design will probably locate these control ports in a different I/O location with different meanings for the bits. There is additional "flash" control logic handled by the PicoPower Redwood chipset. Within the chipset functions, there is advanced memory handling and address range decoder functions that are used to cause the RFA flash memory window to appear in the upper memory block (high memory range) of the system memory. The custom logic mentioned before handles sliding the window to the flash.

The code listings can be found in Appendix E. Due to the volatile nature of software, most likely the listings printed in this Ap-Note will NOT be the most current revision. The full, and current version of the code for this socket services portion of the flash software is available with the ULP Evaluation Kit.

## BOOT ISSUES

As we promised earlier in this document, we will now face the issues involved with booting and how they relate to flash. The ULP Evaluation Board has the ability to boot from floppy disk, hard disk drive, flash RFD, or a PC Card. This necessarily places burdens on both the system as well as the software supporting the flash capabilities of the board (i.e., FTL/SS). The system BIOS, through the use of the BIOS Setup can control what device will be the primary boot device. When flash is the boot device, the following issues need to be handled.

The FTL software must load before boot and typically using the previously discussed INT13h interception method. In order to do this, both the FTL socket services and the FTL itself must load BEFORE the boot is attempted. This is accomplished by storing the software on the system as an option/expansion BIOS. When an x86 PC-like system is starting up, the system BIOS will go out into the UMB and search for code that has the signature (AA55) and proper checksum (zero-sum) of an option/expansion BIOS. The M-System TFFSBIOS code includes utilities to do this in their Integrators TrueFFS FTL package (see the TFFSBIOS Integrator's Guide, TFFSBIOS.DOC, included on the M-Systems disk). Because the ULP uses an Intel Boot Block for the system BIOS, it is a simple matter to store the FTL "Expansion BIOS" or "Option" code in an appropriate spot in memory. For the ULP board, the FTL Socket Services can be found at C800-C8FF. The SS code is approximately 4 KB and is designed to execute in place and as such only requires a small amount of RAM (1K) for its data segment. The FTL itself, in the form of M-System's TFFSBIOS, is located at C900-CFFF. Version 3.2.10 and lower are not specifically designed to execute in place and as such must be copied to DRAM for execution. During the TFFSBIOS code initialization, the code image and data segment are copied to the top of real mode memory (640K memory region) along with the data segment for the socket services. The system's available memory size indicator is appropriately decreased. The combined TFFSBIOS image, buffers,

and data segments take a total of 36 KB-40 KB off of the top of memory. For version 3.2.20 and higher of TFFSBIOS, the code has been designed to operate in place (or execute-in-place) from the flash itself (or a "shadowed" version in Shadow Ram). In this case, only the data segments (including buffers) for TFFSBIOS and the socket services need occupy system ram. The combined requirements total up to approximately 12K taken off the top of memory.

Once the FTL code is recognized and loaded by the BIOS, the flash drive still typically needs to operate as a drive letter that will normally be interrogated by the system BIOS for a boot image. If there are no other "hard drives" installed in the system, this is not a problem as TFFSBIOS automatically grabs the first available "hard drive" spot which just happens to be the usual boot "hard drive." If another hard drive exists in the system, and you want to boot from the FTL flash drive, then special intervention is required. Fortunately, the M-Systems utilities give you an option to force the FTL drive to be drive C: (moving the other hard drives up a spot) and therefore making the FTL flash drive function normally in the boot sequence.

The last major boot issue, (and fairly obvious once you think about it), is the need for a bootable format to exist on the flash from which you are going to boot. The actual requirement is for a master boot record with valid boot information to exist in logical sector 0 of the boot device. This is accomplished in a similar fashion to formatting a hard drive or floppy disk to be bootable. Let's assume the flash in the RFD has never been formatted or has been totally or partially erased. All that is required is to boot up the ULP board normally from another boot device (i.e. a bootable floppy) then invoke the M-System's format utility **TFORMAT /1** (/1 assumes the RFD has been selected as the primary "boot" flash device and is therefore acting as socket 1 out of 2. Use /2 if the RFD is operating secondary to flash in the PC Card socket.) After the FTL format has been laid down on the flash by TFORMAT, it is advisable to reboot the system from the floppy in order for the system to recognize the "new" valid FTL drive. Next, a bootable master boot record has been created on the flash. This is accomplished with the operating system disk **FORMAT** utility, making sure to specify that you wish to transfer the system to the "disk". For DOS, use: **FORMAT /S d:** where *d:* is the drive letter of the FTL flash drive. With a properly formatted and boot image enabled FTL drive as the primary drive, a normal boot can then occur from the flash RFD or Flash Card.



**FTL AVAILABILITY—I WANT IT, WHERE CAN I GET IT?**

“Man, this FTL software sounds GREAT! It will do exactly what I want to do on my system! Where can I get it?” This is a very astute question. There are quite a number of sources for commercially supported FTL. The FTLs from each of the vendors has unique characteristics. Pricing, performance, adaptability, etc. vary from version to version. Please see Appendix A for a listing of the most popular FTL vendors. As stated before, M-Systems is the creator of the specific FTL upon which this document is based.

**CONCLUSION**

The general benefits of flash in an embedded system, and especially a resident flash array operating as a drive with FTL software, are quite obvious. The wide availability and market acceptance of the FTL standard (FTL Is It!) makes the implementation and use of flash in this way a simple matter. And although concepts presented here were presented on a specific platform with a specific set of features, the concepts do apply to virtually any system that needs this kind of capability. With some thought, and a little development work, you too can soon be enjoying the advantages of flash and FTL.

**ADDITIONAL INFORMATION**

**References**

Order Number	Document
290429	28F008SA 8-Mbit (1-Mbit x 8) FlashFile™ Memory Datasheet
272731	Embedded Ultra-Low Power Intel486™ SX Processor Datasheet
272755	Embedded Ultra-Low Power Intel486™ GX Processor Datasheet
292157	AP-605, Implementing a Resident Flash Disk with an Intel386™ EX Embedded Processor
272324	AP-477, Low Voltage Embedded Design
Contact Intel/ Distribution Sales Office	Ultra Low Power Intel486™ SX Evaluation Board Reference Guide/User's Manual

*General Information Hotline*

US/Canada: 1 (800) 628-8686 or (916) 356-7599  
 Japan/APAC: (916) 356-7599  
 Europe: +440 793-69-6776

*Literature Orders:*

US/Canada: 1 (800) 548-4725  
 International: Please contact your local Intel office

BBS: (916) 356-3600 or +4401793-496340  
 FaxBack\*: 1(800) 628-2283 or (916) 356-3105

**Revision History**

Number	Description
-001	Original Version





## APPENDIX A FTL AVAILABILITY

### M-SYSTEMS

*TrueFFS\* (v3.2 and up are FTL)*  
4655 Old Ironsides Dr. Suite #200  
Santa Clara, CA 95054  
(408) 654-5820  
FAX: (408) 654-9107

### System Soft Corporation

*SS FTL*  
313 Speen St.  
Natick, MA 01760  
(508) 651-0088  
FAX: (508) 651-8188

### SCM Microsystem, Inc.

S-FTL (v3.0 and up are FTL)  
131 Albright Way.  
Los Gatos, CA 95030  
(408) 370-4888  
FAX: (408) 370-4880

### Datalight, Inc.

CardTrick FTL  
307 N. Olympic Ave. Suite 200  
Arlington, WA 98223  
(206) 435-8086  
FAX: (206) 435-0253





## APPENDIX B

### PC CARD SOCKET SERVICES

### DESCRIPTIONS AND FUNCTION NUMBERS

*AccessConfigurationSpace* (A2h)

This function is for CardBus. It provides an interface for Card Services to read and write values in the CardBus configuration space.

*AcknowledgeInterrupt* (9Eh)

Returns information about which socket(s) on the adapter specified by the input parameters has had a change in status.

*GetAccessOffsets* (A1h)

Fills the indicated buffer with an array of offsets for adapters using register-based, I/O port access to PC Card memory address space. Used for adapter-specific, low-level, optimized PC Card access routines. Cards that use memory windows, directly mapped into the system memory space do not support this function.

*GetAdapter* (85h)

Returns the current configuration of the specified adapter.

*GetAdapterCount* (80h)

Returns the number of adapters supported by all Socket Services handlers in the host system. Also used to determine if one or more Socket Services handlers are installed.

*GetEDC* (96h)

Returns the current configuration of the EDC (Error Detect and Correct) generator specified by the input parameters.

*GetPage* (8Ah)

Returns the current configuration of the page specified by the input parameters. Only valid for memory windows (WS\_IO is reset for the window).

*GetSetPriorHandler* (9Fh)

Gets, or replaces the entry point of a prior handler for the Adapter specified by the input parameters.

*GetSetSSAddr* (A0h)

Returns code and data area descriptions and provides a way to pass mode-specific data area descriptors to a Socket Services handler.

*GetSocket* (8Dh)

Returns the current configuration of the socket identified by the input parameters.

*GetSSInfo* (83h)

Returns the version / compliance level of the Socket Services interface supporting the adapter specified by the input parameters. It also identifies the adapters serviced by the handler.

*GetStatus* (8Fh)

Returns the current status of the card, socket, controls and indicators for the socket identified by the input parameters.

*GetVendorInfo* (9Dh)

Returns information about the vendor implementing Socket Services for the adapter specified in the input parameters.

*GetWindow* (88h)

Returns the current configuration of the window specified by the input parameters.

<i>InquireAdapter</i>	(84h)	Returns information about the capabilities of the adapter specified by the input parameters.
<i>InquireEDC</i>	(95h)	Returns the capabilities of the EDC generator specified by the input parameters.
<i>InquireSocket</i>	(8Ch)	Returns information about the capabilities of the socket specified by the input parameters.
<i>InquireWindow</i>	(87h)	Returns information about the capabilities of the window specified by the input parameters.
<i>PauseEDC</i>	(99h)	Pauses EDC generation on a configured and in-use EDC generator specified by the input parameters.
<i>ReadEDC</i>	(9ch)	Reads the EDC value computed by the EDC generator specified in the input parameters.
<i>ResetSocket</i>	(90h)	Resets the PC Card in the socket and returns socket hardware to its power-on default state.
<i>ResumeEDC</i>	(9Ah)	Resumes EDC generation on a configured and paused EDC generator specified by the input parameters.
<i>SetAdapter</i>	(86h)	Sets the configuration of the specified adapter.
<i>SetEDC</i>	(97h)	Sets the configuration of the EDC generator specified by the input parameters.
<i>SetPage</i>	(8Bh)	Configures the page specified by the input parameters. Only valid for memory windows. This service is unsupported by PC Card-32.
<i>SetSocket</i>	(8Eh)	Sets the current configuration of the socket identified by the input parameters.
<i>SetWindow</i>	(89h)	Sets the configuration of the window specified by the input parameters.
<i>StartEDC</i>	(98h)	Starts a previously configured EDC generator specified by the input parameters.
<i>StopEDC</i>	(9Bh)	Stops EDC generation on a configured and functioning EDC generator specified by the input parameters.
<i>VendorSpecific</i>	(AEh)	As the name implies, this service is vendor specific. It is reserved for vendors to add proprietary extensions to the Socket Services interface.





## APPENDIX C BIT TWIDDLING

An interesting capability of Intel Flash, in general, is that a flash byte or word can be “written” as many times as you wish *without* having to first erase it and without affecting the component’s lifetime. You might be thinking “Now wait just a minute here! You just said I have erase flash before I write.” That is only necessary only for a certain circumstance: when you need any bits to change from a zero to a one. However, you can *rewrite* flash as long as bits are being changed from 1 to zero, or are not changing at all. The following example illustrates the point:

What	Flash Byte (hex)	Flash Byte (Binary)	OK or Fail
Flash byte Erase:	FF	1 1 1 1 1 1 1 1	
		↓ ↓ ↓ ↓	
Flash byte write to AAh:	AA	1 0 1 0 1 0 1 0	OK!
		↓	
Flash byte write to A8h	A8	1 0 1 0 1 0 0 0	OK!
		↓	
Flash byte write to 28h	28	0 0 1 0 1 0 0 0	OK!
		↓ ↓	
Flash byte write to 00h	00	0 0 0 0 0 0 0 0	OK!
Flash byte write to 01h	0	0 0 0 0 0 0 0 0	FAIL

Erase is required to rewrite this flash byte to anything other than 00h.







## APPENDIX D RELATED THIRD PARTY DOCUMENTATION

### **M-Systems TrueFFS User's Manual**

*M-Systems*  
4655 Old Ironsides Dr. Suite #200  
Santa Clara, CA 95054  
(408) 654-5820  
FAX: (408) 654-9107

### **PT86C768 & PT86C718 "Redwood" Chipset Manual**

*Cirrus Logic / PicoPower*  
3100 West Warren Ave.  
Fremont, CA 94538  
(510) 623-8300  
FAX: (510) 252-6020

### **PC Card Standard**

*PCMCIA*  
2635 North 1st Street  
San Jose, California 95134  
(408) 433-2273  
FAX: (408) 433-9558





## APPENDIX E CODE LISTINGS

```
; $Log: SS.ASM $
; Revision 1.4 1995/09/08 12:55:56 mgianopulos
; InitSocketServices() - Bug fix when booting from PCIC slot.
; Added option to boot from RFA only:InitSocketServices()
; & GetBootDevice().
; Revision 1.1 1995/08/30 09:25:57
; Initial revision
;
; -----
; July/Aug. 1995 Mark Gianopulos - Intel Corp.
; Modified for 486sx(TM) ULP RFA & PCMCIA Socket
; -----
;
; Rev 3.3 28 Feb 1994 10:41:26
; Fixed ResetSocket bug
;
; Rev 3.2 08 Sep 1993 21:18:14
; Version 2.01
;
; Rev 3.1 01 Sep 1993 11:31:06
;
; Rev 3.0 21 Jul 1993 11:38:34
; Speed & IRQ additions
;
; Rev 1.4 25 Jun 1993 12:47:34
; No window, Step B, and other changes
;
; Rev 1.3 30 Mar 1993 10:42:38
; GetVersion change
;
; Rev 1.2 07 Mar 1993 16:58:14
; Power table correction
;
; Rev 1.1 18 Jan 1993 11:44:04
; SS Ver 2.0 final draft
;
; Rev 1.0 12 Jan 1993 16:08:48
; Initial revision.
;

PAGE 78, 132
TITLE SS.ASM
```

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```

COMMENT
/*****
Title:          Intel ULP 486sx RFA Socket Services
                Copyright (c) 1995 Intel Corp.

Author:         Amir Ban
                M-Systems, Ltd.
co-author:     Mark Gianopulos
                Intel Corp.

Copyright (C) by M-Systems Ltd. 1992.  All rights reserved.

*****/

.486

INCLUDE        SSDEFS.INC          ; global definitions
INCLUDE        RFA.INC
INCLUDE        PCIC.INC
INCLUDE        VER.INC             ; version number

                ASSUME CS:NOTHING, DS:NOTHING, ES:NOTHING

ABSO           SEGMENT AT 00H PUBLIC USE16      ; Zero page definition for
vectors
ABSO           ENDS

DGROUP        GROUP _TEXT, _DATA, _INIT

; ---         Subroutine definitions

_TEXT         SEGMENT BYTE PUBLIC 'CODE' USE16

                EXTRN  pGetWindow:NEAR
                EXTRN  pSetWindow:NEAR
                EXTRN  pGetPage:NEAR
                EXTRN  pSetPage:NEAR
                EXTRN  pGetSocket:NEAR
                EXTRN  pSetSocket:NEAR
                EXTRN  pGetStatus:NEAR
                EXTRN  pResetSocket:NEAR

                EXTRN  rGetWindow:NEAR
                EXTRN  rSetWindow:NEAR
                EXTRN  rGetPage:NEAR
                EXTRN  rSetPage:NEAR
                EXTRN  rGetSocket:NEAR
                EXTRN  rSetSocket:NEAR
                EXTRN  rGetStatus:NEAR

_TEXT         ENDS

```

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```

COMMENT
/*****
Customization Definitions
*****/

DISPATCH_ENTRY      STRUC                ; Flash Entry dispatch table
  pFunction          DW      ?            ; Pointer to function
DISPATCH_ENTRY      ENDS

; ----- Equates
NO_OF_PAGES          EQU      1
NO_OF_EDCS           EQU      0          ; no support for EDCs

BOOT_PCIC            EQU      0
BOOT_RFA             EQU      1
BOOT_RFA_ONLY       EQU      2
BOOT_OTHER          EQU      3

_DATA                SEGMENT PARA PUBLIC 'DATA' USE16
; -----
      PUBLIC NoOfAdapters
      PUBLIC NoOfWindows
NoOfAdapters         DB      ?           ; No. of installed adapters
NoOfWindows          DB      ?           ; No. of installed windows
NoOfSockets          DB      ?           ; No. of installed sockets

      PUBLIC RfaSocketNo
      PUBLIC PciSocketNo
RfaSocketNo          DB      0           ; Socket number assigned to RFA
PciSocketNo          DB      0           ; Socket number assigned to PCIC

BootDevice           DB      ?           ; System startup boot device
                                           ; Use BOOT_XXX equates only

      PUBLIC MemWndCaps
      PUBLIC Fastest
WindowCharsTable     LABEL NEAR
MemWndCaps           DW      0000000011001111b ; Capabilities (see SS document-
InquireWindow)
MinWinBase           DW      ?           ; Minimum base address (4K blocks)
MaxWinBase           DW      ?           ; Maximum base address (4K blocks)
MinWinSize           DW      4           ; 16K, Minimum window size (4K
blocks)
MaxWinSize           DW      4           ; 16K, Maximum window size (4K
blocks)
ReqGran              DW      1           ; 4K, Window size granularity
ReqBase              DW      4           ; 16K, Base address alignment
CardOffAlign         DW      1           ; Card offset alignment
Slowest              DB      4           ; 100ns, Slowest access speed
Fastest              DB      4           ; 100ns, Fastest access speed
WindowTableLength = $ - WindowCharsTable

RoutineAddress       DW      ?

_DATA                ENDS

```

```

__TEXT      SEGMENT BYTE PUBLIC 'CODE' USE16

; -----      Flash Entry point dispatch table, contains pointer to
function
;              and whether media check needs to be performed.

SSDispatch  DISPATCH_ENTRY    <GetAdapterCount>      ; 80
            DISPATCH_ENTRY    <NotImplemented>       ; 81
            DISPATCH_ENTRY    <NotImplemented>       ; 82
            DISPATCH_ENTRY    <GetSSInfo>            ; 83
            DISPATCH_ENTRY    <InquireAdapter>       ; 84
            DISPATCH_ENTRY    <GetAdapter>           ; 85
            DISPATCH_ENTRY    <SetAdapter>           ; 86
            DISPATCH_ENTRY    <InquireWindow>        ; 87
            DISPATCH_ENTRY    <GetWindow>            ; 88
            DISPATCH_ENTRY    <SetWindow>            ; 89
            DISPATCH_ENTRY    <GetPage>              ; 8a
            DISPATCH_ENTRY    <SetPage>              ; 8b
            DISPATCH_ENTRY    <InquireSocket>        ; 8c
            DISPATCH_ENTRY    <GetSocket>            ; 8d
            DISPATCH_ENTRY    <SetSocket>            ; 8e
            DISPATCH_ENTRY    <GetStatus>            ; 8f
            DISPATCH_ENTRY    <ResetSocket>          ; 90
            DISPATCH_ENTRY    <NotImplemented>       ; 91
            DISPATCH_ENTRY    <NotImplemented>       ; 92
            DISPATCH_ENTRY    <NotImplemented>       ; 93
            DISPATCH_ENTRY    <NotImplemented>       ; 94
            DISPATCH_ENTRY    <NotImplemented>       ; 95
            DISPATCH_ENTRY    <NotImplemented>       ; 96
            DISPATCH_ENTRY    <NotImplemented>       ; 97
            DISPATCH_ENTRY    <NotImplemented>       ; 98
            DISPATCH_ENTRY    <NotImplemented>       ; 99
            DISPATCH_ENTRY    <NotImplemented>       ; 9a
            DISPATCH_ENTRY    <NotImplemented>       ; 9b
            DISPATCH_ENTRY    <NotImplemented>       ; 9c
            DISPATCH_ENTRY    <GetVendorInfo>        ; 9d
            DISPATCH_ENTRY    <NotImplemented>       ; 9e
            DISPATCH_ENTRY    <GetSetPriorHandler>   ; 9f
            DISPATCH_ENTRY    <GetSetSSAddr>        ; a0

NUM_Functions = ($ - OFFSET SSDispatch) / TYPE DISPATCH_ENTRY

Implementor  DB      'HEB_SS  '

SocketCharsTable LABEL NEAR
                DW      1                ; Interface-type support (memory)
                DD      0                ; IRQ levels inverting IREQ line
                DD      0                ; IRQ levels not inverting IREQ
line

SocketTableLength = $ - SocketCharsTable

```

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```

AdapterCharsTable LABEL NEAR
    DW 0 ; Capabilities (individual)
    DW 0 ; Active high IRQ levels
    DW 08h ; - required by SysSoft Card
Services DD 0 ; Active low IRQ levels

PowerManagementTable LABEL NEAR
    DW 0 ; No entries - assume required pwr
exists

AdapterTableLength = $ - AdapterCharsTable

VendorInfo LABEL NEAR
    DB 'Intel ULP Eval Board SS', 0
VendorInfoLength = $ - VendorInfo

PAGE
COMMENT
/*****

Procedure: GetAdapterCount

Entry: No significant registers

Exit: [AL] = Number of adapters supported
      [CX] = 'SS' (Socket Services id)

*****/

GetAdapterCount PROC NEAR
    ASSUME CS:_TEXT, DS:DGROUP, ES:NOTHING

    MOV AL, NoOfAdapters ; One adapter
    MOV CX, 'SS'

    CLC
    RET

GetAdapterCount ENDP
    
```

```

                PAGE
COMMENT
/*****
Procedure:  GetSSInfo

Entry:      [AL] = Adapter

Exit:       [AX] = 0 (Version 1.0 compatibility)
            [BX] = Socket Services version number
            [CH] = Number of adapters supported by handler
            [CL] = First adapter supported by handler

*****/

GetSSInfo  PROC NEAR
            ASSUME     CS:_TEXT, DS:DGROUP, ES:NOTHING

            MOV       BX, PCMCIA_COMPLIANCE    ; PCMCIA SS Release
            MOV       CH, 1                    ; This SS supports only One
adapter
            MOV       CL, NoOfAdapters        ; first adapter number
            DEC       CL                       ; = total system adapters - 1

            XOR       AX, AX                   ; Return success (resets [CF])
            RET

GetSSInfo  ENDP

```

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```

                PAGE
COMMENT
/*****
Procedure:  InquireAdapter

Entry:      [AL] = Adapter
            [ES]:[DI] = Pointer to buffer for adapter
characteristics
            and power management tables

Exit:       [BH] = Number of windows
            [BL] = Number of sockets
            [CX] = Number of EDCs
            [ES]:[DI] = Unchanged

*****/

InquireAdapter PROC    NEAR
                ASSUME     CS:_TEXT, DS:DGROUP, ES:NOTHING

                MOV     CX, AdapterTableLength
                MOV     ES:[DI + 2], CX           ; return wDataLength
                CMP     CX, ES:[DI]             ; do we need to truncate?
                JB      IA2
                MOV     CX, ES:[DI]           ; yes, truncate returned

data
IA2:

                PUSH   SI
                PUSH   DI
                MOV    SI, OFFSET DGROUP:AdapterCharsTable
                ADD    DI, 4
                REP   MOVSB                   ; copy AdapterCharsTable to ES:DI
                POP    DI
                POP    SI

                MOV    BH, NoOfWindows
                MOV    BL, NoOfSockets
                MOV    CX, NO_OF_EDCS

                XOR    AH, AH                 ; Return success (resets [CF])
                RET

InquireAdapter ENDP

```

292173-17

```

                PAGE
COMMENT
/*****

    Procedure:  GetAdapter

    Entry:      [AL] = Adapter

    Exit:       [DH] = Adapter State:
                Bit 0   Reduced power consumption
                Bit 1   State information preserved
                [DI] = IRQ level used for status change interrupts

    *****/

GetAdapter  PROC  NEAR
            ASSUME  CS:_TEXT, DS:DGROUP, ES:NOTHING

            XOR    DH, DH
            XOR    DI, DI

            XOR    AH, AH           ; Return success (resets [CF])
            RET

GetAdapter  ENDP

                PAGE
COMMENT
/*****

    Procedure:  SetAdapter

    Entry:      [AL] = Adapter
                [DH] = Adapter State:
                Bit 0   Reduced power consumption
                Bit 1   State information preserved
                [DI] = IRQ level used for status change interrupts

    Exit:       All registers preserved

    *****/

SetAdapter  PROC  NEAR
            ASSUME  CS:_TEXT, DS:DGROUP, ES:NOTHING

            XOR    AH, AH           ; Return success (resets [CF])
            RET

SetAdapter  ENDP

```

292173-18

```

PAGE
COMMENT
~*****
Procedure:  InquireWindow

Entry:      [AL] = Adapter
            [BH] = Window
            [ES]:[DI] = Pointer to buffer for window characteristics

Exit:       [BL] = Capabilities:
            Bit 0  Common memory
            Bit 1  Attribute memory
            Bit 2  I/O space
            Bit 7  Wait supported
            [CX] = Bit-map of assignable sockets
            [ES]:[DI] = buffer filled with window characteristics
~*****~

InquireWindow PROC  NEAR
                ASSUME      CS:_TEXT, DS:DGROUP, ES:NOTHING

                CMP        BH, NoOfWindows
                JB         IW2
                MOV        AH, BadWindow
                STC
                RET

IW2:

                CMP        BH, PCIC_WINDOW
                JNE        IW3

                ; PCIC Window Characteristics (Window 1, Socket ?)
                MOV        MinWinBase, pMinWinBase
                MOV        MaxWinBase, pMaxWinBase

                MOV        BL, 3                ; Common & Attribute memory
                XOR        CX, CX
                MOV        CL, PccSocketNo     ; window available for this socket
                INC        CX                ; CX = PccSocketNo + 1
                JMP        IW4

IW3:

                ; RFA Window Characteristics (Window 0, Socket ?)
                MOV        MinWinBase, rMinWinBase
                MOV        MaxWinBase, rMaxWinBase

                MOV        BL, 1                ; Common memory only
                XOR        CX, CX
                MOV        CL, RfaSocketNo     ; window available for this socket
                INC        CX                ; CX = RfaSocketNo + 1

IW4:

                PUSH        CX
                MOV        CX, WindowTableLength
                MOV        ES:[DI + 2], CX
                CMP        CX, ES:[DI]
                JB         IW5
                MOV        CX, ES:[DI]

```

```

IW5:
        PUSH    SI
        PUSH    DI
        MOV     SI, OFFSET DGROUP:WindowCharsTable
        ADD     DI, 4
        REP    MOVSB
        POP     DI
        POP     SI
        POP     CX

        XOR     AH, AH                ; Return success (resets [CF])
        RET

```

```
InquireWindow ENDP
```

```
        PAGE
```

```
COMMENT
```

```
~*****~
```

```
        Procedure:  GetWindow
```

```
        Entry:      [AL] = Adapter
                   [BH] = Window
```

```
        Exit:       [BL] = Socket
                   [CX] = Window size
                   [DH] = Window state:
                       Bit 0  Memory or I/O
                       Bit 1  Disabled or enabled
                       Bit 2  8-bit path or 16-bit path
                       Bit 3  Subdivided into pages
                   [DL] = Access speed
                   [DI] = Window base address

```

```
*****~
```

```
GetWindow  PROC  NEAR
            ASSUME  CS:_TEXT, DS:DGROUP, ES:NOTHING

```

```
            CMP    BH, NoOfWindows
            JB     GW2
            MOV    AH, BadWindow
            STC
            RET

```

```
GW2:
            CMP    BH, PCIC_WINDOW
            JNE    GW3

```

```
            CALL   pGetWindow
            JMP    GW4

```

```
GW3:
            CALL   rGetWindow

```

```
GW4:
            RET

```

```
GetWindow  ENDP
```

292173-20

```

                PAGE
COMMENT
~*****~
Procedure:  SetWindow

Entry:      [AL] = Adapter
            [BH] = Window
            [BL] = Socket
            [CX] = Window size
            [DH] = Window state:
                Bit 0  Memory or I/O
                Bit 1  Disabled or enabled
                Bit 2  8-bit path or 16-bit path
                Bit 3  Subdivided into pages/EISA I/O mapping
                Bit 4  EISA common I/O accesses
            [DL] = Access speed
            [DI] = Window base address

Exit:      All registers preserved

*****~

SetWindow  PROC  NEAR
            ASSUME  CS:_TEXT, DS:DGROUP, ES:NOTHING

            CMP     BH, NoOfWindows
            JB      SW2
            MOV     AH, BadWindow
            STC
            RET

SW2:
            ; Check for valid socket
            CMP     BL, NoOfSockets
            JB      SW3
            MOV     AH, BadSocket
            STC
            RET

SW3:
            CMP     BH, PCIC_WINDOW
            JNE     SW4

            CALL    pSetWindow
            JMP     SW5

SW4:
            CALL    rSetWindow

SW5:
            RET

SetWindow  ENDP
    
```

292173-21



```

                PAGE
COMMENT
~*****~
Procedure:  GetPage

Entry:      [AL] = Adapter
            [BH] = Window
            [BL] = Page

Exit:       [DL] = Page state
            Bit 0 Common space or attribute space
            Bit 1 Disabled or enabled
            [DI] = Memory card offset (4 KByte units)

*****~

GetPage      PROC   NEAR
              ASSUME      CS:_TEXT, DS:DGROUP, ES:NOTHING

              CMP    BH, NoOfWindows
              JB     GP2
              MOV    AH, BadWindow
              STC
              RET

GP2:         CMP    BL, NO_OF_PAGES
              JB     GP3
              MOV    AH, BadPage
              STC
              RET

GP3:         CMP    BH, PCIC_WINDOW
              JNE    GP4

              CALL   pGetPage
              JMP    GP5

GP4:         CALL   rGetPage

GP5:         RET

GetPage      ENDP
    
```

292173-22





```

                PAGE
COMMENT
~*****
Procedure:    SetPage

Entry:        [AL] = Adapter
               [BH] = Window
               [BL] = Page
               [DL] = Page state
                 Bit 0 Common space or attribute space
                 Bit 1 Disabled or enabled
                 Bit 2 WP - Write protection
               [DI] = Memory card offset (4 KByte units)

Exit:         All registers preserved

*****~

SetPage      PROC    NEAR
              ASSUME    CS:_TEXT, DS:DGROUP, ES:NOTHING

              CMP      BH, NoOfWindows
              JB       SP2
              MOV      AH, BadWindow
              STC
              RET

SP2:         CMP      BL, NO_OF_PAGES
              JB       SP3
              MOV      AH, BadPage
              STC
              RET

SP3:         CMP      BH, PCIC_WINDOW
              JNE      SP4

              CALL    pSetPage
              JMP     SP5

SP4:         CALL    rSetPage

SP5:         RET

SetPage      ENDP
    
```

292173-23

```

COMMENT
~*****
Procedure:  InquireSocket

Entry:      [AL] = Adapter
            [BL] = Socket
            [ES]:[DI] = Pointer to Socket Characteristics

Exit:      [BH] = Status Change Interrupt Capabilities
            Bit 0  Write Protect Change
            Bit 1  Card Lock Change
            Bit 2  Ejection Request
            Bit 3  Insertion request
            Bit 4  Battery dead change
            Bit 5  Battery warning change
            Bit 6  Ready change
            Bit 7  Card detect change
            [DH] = Status Change reporting capabilities:
            Bit 0  Write Protect Change
            Bit 1  Card Lock Change
            Bit 2  Ejection Request
            Bit 3  Insertion request
            Bit 4  Battery dead change
            Bit 5  Battery warning change
            Bit 6  Ready change
            Bit 7  Card detect change
            [DL] = Control and indicators capabilities
            Bit 0  Write protect status
            Bit 1  Card lock status
            Bit 2  Motorized card ejection
            Bit 3  Motorized card ejection
            Bit 4  Card lock
            Bit 5  Battery status
            Bit 6  Busy status
            Bit 7  XIP status
            [ES]:[DI] = Unchanged

*****~

InquireSocket PROC    NEAR
                   ASSUME     CS:_TEXT, DS:DGROUP, ES:NOTHING

                   CMP        BL, NoOfSockets
                   JB         IS2
                   MOV        AH, BadSocket
                   STC
                   RET

IS2:
                   PUSH       CX
                   MOV        CX, SocketTableLength
                   MOV        ES:[DI + 2], CX
                   CMP        CX, ES:[DI]
                   JB         IS3
                   MOV        CX, ES:[DI]

IS3:
                   PUSH       SI

```



```
PUSH    DI
MOV     SI, OFFSET DGROUP:SocketCharsTable
ADD     DI, 4
REP     MOVSB
POP     DI
POP     SI
POP     CX

MOV     BH, 0
MOV     DH, 80h      ; Report card detect (always inserted)
MOV     DL, 0

XOR     AH, AH      ; Return success (resets [CF])
RET
```

```
InquireSocket ENDP
```

292173-25

```

~*****
Procedure:  GetSocket

Entry:      [AL] = Adapter
           [BL] = Socket

Exit:       [BH] = Status change interrupt mask
           Bit 0  Write Protect Change
           Bit 1  Card Lock Change
           Bit 2  Ejection Request
           Bit 3  Insertion request
           Bit 4  Battery dead change
           Bit 5  Battery warning change
           Bit 6  Ready change
           Bit 7  Card detect change
           [CH] = VCC level
           [CL] = VPP1 & VPP2 levels
           [DH] = Socket state
           Bit 0  Write Protect Change
           Bit 1  Card Lock Change
           Bit 2  Ejection Request
           Bit 3  Insertion request
           Bit 4  Battery dead change
           Bit 5  Battery warning change
           Bit 6  Ready change
           Bit 7  Card detect change
           [DL] = Control and indicators state
           Bit 0  Write protect status
           Bit 1  Card lock status
           Bit 2  Motorized card ejection
           Bit 3  Motorized card ejection
           Bit 4  Card lock
           Bit 5  Battery status
           Bit 6  Busy status
           Bit 7  XIP status
           [DI] = IRQ level steering/Interface type

*****~
GetSocket  PROC  NEAR
           ASSUME      CS:_TEXT, DS:DGROUP, ES:NOTHING

           CMP  BL, NoOfSockets
           JB  GS2
           MOV  AH, BadSocket
           STC
           RET

GS2:
           CMP  BL, PcicSocketNo
           JNE  GS3

           CALL pGetSocket
           JMP  GS4

GS3:
           CALL rGetSocket

GS4:
           RET
GetSocket  ENDP

```

292173-26



```

~*****
Procedure:  SetSocket

Entry:      [AL] = Adapter
            [BL] = Socket

Exit:       [BH] = Status change interrupt mask
            Bit 0  Write Protect Change
            Bit 1  Card Lock Change
            Bit 2  Ejection Request
            Bit 3  Insertion request
            Bit 4  Battery dead change
            Bit 5  Battery warning change
            Bit 6  Ready change
            Bit 7  Card detect change

            [CH] = VCC level
            [CL] = VPP1 & VPP2 levels
            [DH] = Socket state
            Bit 0  Write Protect Change
            Bit 1  Card Lock Change
            Bit 2  Ejection Request
            Bit 3  Insertion request
            Bit 4  Battery dead change
            Bit 5  Battery warning change
            Bit 6  Ready change
            Bit 7  Card detect change

            [DL] = Control and indicators state
            Bit 0  Write protect status
            Bit 1  Card lock status
            Bit 2  Motorized card ejection
            Bit 3  Motorized card ejection
            Bit 4  Card lock
            Bit 5  Battery status
            Bit 6  Busy status
            Bit 7  XIP status

            [DI] = IRQ level steering/Interface type

Exit:       All register preserved
*****~
SetSocket   PROC   NEAR
            ASSUME     CS:_TEXT, DS:DGROUP, ES:NOTHING

            CMP     BL, NoOfSockets
            JB      SS2
            MOV     AH, BadSocket
            STC
            RET

SS2:
            CMP     BL, PciSocketNo
            JNE     SS3

            CALL    pSetSocket
            JMP     SS4

SS3:
            CALL    rSetSocket

SS4:
            RET

SetSocket   ENDP

```

```

~*****
Procedure:  GetStatus

Entry:      [AL] = Adapter
            [BL] = Socket

Exit:       [BH] = Current card state:
            Bit 0  Write protect
            Bit 1  Card Locked
            Bit 2  Ejection Request
            Bit 3  Insertion Request
            Bit 4  Battery dead
            Bit 5  Battery low
            Bit 6  Card ready
            Bit 7  Card detected
            [DH] = Socket state
            [DL] = Control and indicator state
            [DI] = IRQ level steering/Interface type

*****~

GetStatus  PROC  NEAR
            ASSUME  CS:_TEXT, DS:DGROUP, ES:NOTHING

            CMP    BL, NoOfSockets
            JB     GC2
            MOV    AH, BadSocket
            STC
            RET

GC2:
            CMP    BL, PciSocketNo
            JNE   GC3

            CALL  pGetStatus
            JMP   GC4

GC3:
            CALL  rGetStatus

GC4:
            RET

GetStatus  ENDP
    
```

292173-28



```

~*****
  Procedure:  ResetSocket
    Entry:    [AL] = Adapter
             [BL] = Socket

    Exit:     All registers preserved
~*****~
ResetSocket PROC NEAR
    ASSUME    CS:_TEXT, DS:DGROUP, ES:NOTHING
    CMP      BL, NoOfSockets
    JB      RS2
    MOV      AH, BadSocket
    STC
    RET

RS2:
    CMP      BL, PciSocketNo
    JNE     RS3
    CALL    pResetSocket
    JMP     RS4

RS3:
    XOR      AH, AH                ; Return success (resets [CF])

RS4:
    RET
ResetSocket ENDP

~*****
  Procedure:  GetVendorInfo
    Entry:    [AL] = Adapter
             [BL] = Type
             [ES]:[DI] = Vendor info buffer

    Exit:     [DX] = Vendor Release
~*****~
GetVendorInfo PROC NEAR
    ASSUME    CS:_TEXT, DS:DGROUP, ES:NOTHING
    CMP      BL, 0
    JNE     GV1
    PUSH     CX
    MOV      CX, VendorInfoLength
    MOV      ES:[DI + 2], CX
    CMP      CX, ES:[DI]
    JB      GV3
    MOV      CX, ES:[DI]

GV3:
    PUSH     SI
    PUSH     DI
    MOV      SI, OFFSET DGROUP:VendorInfo
    ADD      DI, 4
    REP     MOVSB
    POP      DI
    POP      SI
    POP      CX

GV1:
    MOV      DX, VENDOR_VERSION ; Our version no.
    XOR      AH, AH                ; Return success (resets [CF])
    RET
GetVendorInfo ENDP

```

292173-29

```

/*****
Unimplemented Socket Services
*****/

NotImplemented PROC NEAR
    ASSUME     CS:_TEXT, DS:NOTHING, ES:NOTHING

    MOV     AX, UnsupportedFunction
    STC

    RET
NotImplemented ENDP

/*****
Procedure:  _SocketServices
Purpose:    Main entry point of Socket Services
Entry:     [AH] = Function, starting at 80h
           Others vary by function
Exit:      [CF] = Status
           If [CF] set, [AX] = error
           Others vary by function
*****/
PUBLIC _SocketServices

_SocketServices PROC FAR
    ASSUME     CS:_TEXT, DS:DGROUP, ES:NOTHING

    AND     AH, 7Fh
    CMP     AH, NUM_Functions    ; Valid request ?
    JB     SoS2                  ; Yes, continue

    MOV     AX, UnsupportedFunction ; No, invalid command
    STC
    RET
    ; Set carry bit to indicate error

SoS2:
    PUSH    AX
    PUSH    BX
    XOR     AL, AL
    XCHG   AH, AL
    ADD     AX, AX
    MOV     BX, AX
    ADD     BX, OFFSET DGROUP:SSDispatch ; Add base address of tbl
    MOV     AX, CS:[BX.pFunction]
    MOV     RoutineAddress, AX
    POP     BX
    POP     AX
    CALL    [RoutineAddress]    ; Perform requested function
    ; AX contains error code from function

    RET
_SocketServices ENDP
_TEXT     ENDS

```

292173-30



```

_INIT          SEGMENT BYTE PUBLIC 'INIT' USE16
               ASSUME          CS:DGROUP, DS:DGROUP, ES:NOTHING

               EXTRN   PpicInit:NEAR
               EXTRN   RfaInit:NEAR

szSignonMsg    DB      CR, LF
               DB      'Intel ULP 486sx (TM) Evaluation Board PCMCIA '
               DB      sPCMCIA_COMPLIANCE, ' Socket Services Interface.', CR,
LF
               DB      '(C) Copyright 1992, M-Systems Ltd.', CR, LF
               DB      '(C) Copyright 1995, Intel Corporation, Version '
               DB      sVENDOR_VERSION, CR, LF
               DB      LF
               DB      EOS

szAbortMsg     DB      CR, LF
               DB      'Initialization aborted.', CR, LF
               DB      EOS

szPressKeyMsg  DB      CR, LF
               DB      'Press any key to continue ... ', CR, LF
               DB      EOS

szLongDelay    DB      'Hard Disk Drive detected!', CR, LF
               DB      'This will cause a 45 second delay.  Please wait...'
               DB      CR, LF
               DB      CR, LF
               DB      EOS

szSelectBootDevice DB    ' Select Boot Device', CR, LF, LF
               DB      ' 1) PCMCIA Slot + RFA', CR, LF
               DB      ' 2) RFA + PCMCIA Slot', CR, LF
               DB      ' 3) RFA only', CR, LF
               DB      CR, LF
               DB      ' Enter a number: '
               DB      EOS

szResponse1    DB      '1'
               DB      CR, LF
               DB      EOS

szResponse2    DB      '2'
               DB      CR, LF
               DB      EOS

szResponse3    DB      '3'
               DB      CR, LF
               DB      EOS

szCrLf         DB      CR, LF
               DB      EOS

```

292173-31

```

/*****
Procedure:  _InitSocketServices
Purpose:    Do first-time initialization

Entry:     [AL] = Adapter no.
           [DX] = BIOS Expansion segment (0 = not applicable)
           [ES]:[DI] = 8-Character implementation name buffer

Exit:      [AX] = status

Locate Adapter and establish Adapter Data
Install Socket Services

*****/

        PUBLIC _InitSocketServices
_InitSocketServices PROC NEAR
        ASSUME     CS:DGROUP, DS:DGROUP, ES:NOTHING

        PUSH     DX

        INC      AL
        MOV      NoOfAdapters, AL

        MOV      DX, OFFSET DGROUP:szSignonMsg
        CALL     PrintMsg

        PROTO1_BUGFIX

        CALL     _GetBootDevice

        CMP      AL, BOOT_PCIC
        JNE      Init1

        MOV      PciSocketNo, 0
        MOV      RfaSocketNo, 1
        MOV      NoOfSockets, 2
        MOV      NoOfWindows, 2
        JMP      Init2

Init1:   MOV      RfaSocketNo, 0
        MOV      PciSocketNo, 1
        MOV      NoOfSockets, 2
        MOV      NoOfWindows, 2

        CMP      AL, BOOT_RFA_ONLY
        JNE      Init2
        MOV      PciSocketNo, -1
        MOV      NoOfSockets, 1
        MOV      NoOfWindows, 1
        JMP      Init3

Init2:   ; (3) Initialize Socket Services
        CALL     PciInit
        JNC      Init3

```

292173-32

```

        MOV     DX, OFFSET DGROUP:szAbortMsg
        CALL   PrintMsg

        MOV     DX, OFFSET DGROUP:szPressKeyMsg
        CALL   PressKey          ; Wait for acknowledgement

        POP     DX
        RET

Init3:
        CALL   RfaInit
        POP     DX

        ; Copy Implementor string into ES:DI buffer
        PUSH   SI
        PUSH   CX

        MOV     CX, 8
        MOV     SI, OFFSET DGROUP:Implementor
        REP    MOVSB

        POP     CX
        POP     SI

        XOR     AX, AX          ; Return success (resets [CF])
        RET

_InitSocketServices ENDP

COMMENT
/*****
Procedure:  _GetBootDevice
Purpose:    Determine which device is primary boot device.
            (i.e. IDE hard drive, RFA, or PCMCIA slot)
Exit:      [AL] = BOOT_xxx (boot device)
*****/

        PUBLIC _GetBootDevice
_GetBootDevice PROC NEAR
        ASSUME CS:DGROUP, DS:DGROUP, ES:NOTHING

        ; Does IDE Hard drive exist?
        PUSH   ES

        MOV     AX, ABSO          ; Point to ABSO segment
        MOV     ES, AX
        MOV     AL, ES:HardDrive

        POP     ES

        CMP     AL, 0
        JE     GBD2
        ; Yes!

```

```

        MOV     DX, OFFSET DGROUP:szLongDelay
        CALL   PrintMsg
        MOV     AL, BOOT_OTHER
        JMP     GBD3 ; leave

GBD2:
        ; No! Hard drive does not exist, therefore
        ; prompt for boot device (RFA or PCMCIA Slot)
        MOV     DX, OFFSET DGROUP:szSelectBootDevice
        CALL   PrintMsg

GBD6:
        XOR     AH, AH
        INT     16h ; Read character from keyboard buffer

        OR     AL, AL
        JZ     GBD5
        MOV     DX, OFFSET DGROUP:szResponse1 ; Echo character
        CMP     AL, 31h
        JE     GBD7
        MOV     DX, OFFSET DGROUP:szResponse2
        CMP     AL, 32h
        JE     GBD7
        MOV     DX, OFFSET DGROUP:szResponse3

GBD7:
        CALL   PrintMsg
        MOV     DX, OFFSET DGROUP:szCrLf
        CALL   PrintMsg

        AND     AX, 000Fh
        DEC     AL
        CMP     AL, BOOT_OTHER
        JAE    GBD6
        JMP     GBD3

GBD5:
        ; Send Function key back to keyboard buffer for BIOS to process
        CMP     AH, 3Ch ; F2 key?
        JNE    GBD6
        MOV     AH, 05
        XOR     CL, CL
        MOV     CH, 3Ch
        INT     16h

GBD3:
        RET

_GetBootDevice ENDP

_INIT      ENDS

        END

```

292173-34

```

; $Log: RFA.ASM $
; Revision 1.1 1995/08/30 09:25:06
; Initial revision
;
; -----
; July/Aug. 1995          Mark Gianopulos - Intel Corp.
; Created for 486sx(TM) ULP RFA & PCMCIA Socket
; -----
;

        PAGE    78, 132
        TITLE   RFA.ASM
COMMENT ~*****

        Title:      RFA.ASM

        Purpose:    RFA Specific functions for ULP Socket Services

        Author:     Mark Gianopulos
                   Intel Corporation

        Copyright (C) by Intel 1995. All rights reserved.

*****~

.486

INCLUDE    SSDEFS.INC          ; global definitions
INCLUDE    RFA.INC

        ASSUME    CS:NOTHING, DS:NOTHING, ES:NOTHING

DGROUP    GROUP    _TEXT, _DATA, _INIT

_DATA     SEGMENT PARA PUBLIC 'DATA' USE16

EXTRN     RfaSocketNo:BYTE          ; ss.asm
EXTRN     Fastest:BYTE              ; ss.asm

MemCardOffset    DW    ?                ; Current card offset (can't read
                                           ; page register)

        DATA    ENDS

```

```

_TEXT          SEGMENT BYTE PUBLIC 'CODE' USE16
; -----
; Proc: SetWindowBaseAddr
; Desc: Sets the RFA system window
;       to the value specified in BX
; I: BX - contains C8, D0, D4, D8, or DC
;
; Trashes: AX, BX, DX
; -----
SetWindowBaseAddr  PROC   NEAR

                SHR     BX, 2                ; D0 ==> 34 (bits14-bits21)
                MOV     BH, 80h             ; enable GPCS2#, (bits22-bits23)

                CLI                                     ; Must not be interrupted
                IODELAY
                MOV     DX, REDWOOD_INDEX   ; Set base address
                MOV     AX, WINDOW_BASE_REG
                OUT     DX, AX
                IODELAY
                MOV     DX, REDWOOD_DATA
                MOV     AX, BX
                OUT     DX, AX

                IODELAY
                MOV     DX, REDWOOD_INDEX   ; Enable base address
                MOV     AX, WINDOW_BASE_EN_REG
                OUT     DX, AX
                IODELAY
                MOV     DX, REDWOOD_DATA
;                MOV     AX, 3FFCh           ; assumes 64K windows only!
;                MOV     AX, 3FFEh           ; assumes 32K windows only!
                MOV     AX, 3FFFh           ; assumes 16K windows only!
                OUT     DX, AX
                STI

                RET

SetWindowBaseAddr  ENDP

```

292173-36

```

; -----
; Proc: GetWindowBaseAddr
; Desc: Returns the RFA system window
;       address in 4K units. (ie: C8, D0, D4, D8, or DC)
; O: AX - returns address in AX register
;
; Trashes: AX, DX
; -----
GetWindowBaseAddr PROC NEAR
    ASSUME CS:_TEXT, DS:DGROUP, ES:NOTHING

    CLI ; Must not be interrupted
    IODELAY
    MOV DX, REDWOOD_INDEX ; Get base address
    MOV AX, WINDOW_BASE_REG
    OUT DX, AX
    IODELAY
    IN AX, REDWOOD_DATA
    STI

    AND AX, 03FFh ; mask out top 6 bits
    SHL AX, 2 ; 34 ==> D0

    RET
GetWindowBaseAddr ENDP

COMMENT ~*****~
Procedure: rGetWindow

Entry: [AL] = Adapter
       [BH] = Window

Exit: [BL] = Socket
      [CX] = Window size
      [DH] = Window state:
          Bit 0 Memory or I/O
          Bit 1 Disabled or enabled
          Bit 2 8-bit path or 16-bit path
          Bit 3 Subdivided into pages
      [DL] = Access speed
      [DI] = Window base address
*****~

PUBLIC rGetWindow
rGetWindow PROC NEAR
    ASSUME CS:_TEXT, DS:DGROUP, ES:NOTHING

    ;[BL] = Socket
    MOV BL, RfaSocketNo

    ;[DI] = Window base address
    call GetWindowBaseAddr ; returned in AX as 4K unit
    MOV DI, AX

    ;[CX] = Window size
    MOV DX, WINDOW_SIZE_REG
    IN AL, DX
    PUSH AX ; save register contents
    
```

292173-37

```

    NOT    AL                ; complement
    AND    AL, 03h          ; mask off all but bits 0 and 1
    INC    AL                ; add 1
    MOV    CL, AL           ; save in cl (shift factor)
    MOV    AX, 01           ; bit to be shifted (2,4,8,16)
    SHL    AL, CL
    MOV    CX, AX

    ;[DH] = Window state
    POP    AX                ; restore register
    MOV    DH, 0            ; 8bit + disabled + mem
    TEST   AL, 04h
    JZ     rGW6
    OR     DH, 10b          ; 8bit + enabled + mem

rGW6:
    ;[DL] = Access speed
    MOV    DL, Fastest

    XOR    AH, AH           ; Return success (resets [CF])
    RET

rGetWindow ENDP

COMMENT ~*****~
Procedure:  rSetWindow

Entry:      [AL] = Adapter
            [BH] = Window
            [BL] = Socket
            [CX] = Window size
            [DH] = Window state:
                Bit 0  Memory or I/O
                Bit 1  Disabled or enabled
                Bit 2  8-bit path or 16-bit path
                Bit 3  Subdivided into pages/EISA I/O mapping
                Bit 4  EISA common I/O accesses
            [DL] = Access speed
            [DI] = Window base address

Exit:       All registers preserved
*****~

PUBLIC rSetWindow
rSetWindow PROC NEAR
    ASSUME     CS:_TEXT, DS:DGROUP, ES:NOTHING

    CMP    BL, RfaSocketNo
    JE     rSW1
    MOV    AH, BadSocket
    STC
    RET

rSW1:
    ; Check for valid window size
    MOV    AH, WINDOW_SIZE_08K
    CMP    CX, 2                ; 8K
    JE     rSW4
    MOV    AH, WINDOW_SIZE_16K

```

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```

    CMP    CX, 4                ; 16K
    JE     rSW4
    MOV    AH, WINDOW_SIZE_32K
    CMP    CX, 8                ; 32k
    JE     rSW4
    MOV    AH, WINDOW_SIZE_64K
    CMP    CX, 10h             ; 64K
    JE     rSW4
    CMP    CX, 0                ; 64K - MAX(WORD)
    JE     rSW4
    MOV    AH, BadSize
    STC
    RET

rSW4:
    ; Set new window size
    PUSH  DX
    MOV   DX, WINDOW_SIZE_REG
    IN    AL, DX
    AND   AL, 0FCh             ; Clear bits 0,1
    OR    AL, AH               ; assign bits 0,1
    OUT   DX, AL
    POP   DX

    ; Check for valid Window Base address
    CMP   DI, rMinWinBase
    JB    rSW5
    CMP   DI, rMaxWinBase
    JBE   rSW6

rSW5:
    MOV   AH, BadOffset
    STC
    RET

rSW6:
    ; Set new window base address
    PUSH  DX
    MOV   BX, DI
    call  SetWindowBaseAddr
    POP   DX

    ; Check for valid window state
    CMP   DH, 0                ; Memory, disabled, 8bit, 1page
    JE    rSW7
    CMP   DH, 2                ; Memory, enabled, 8bit, 1page
    JE    rSW8
    MOV   AH, BadAttribute
    STC
    RET

rSW7:
    ; disable window
    PUSH  DX
    MOV   DX, WINDOW_ENABLE_REG
    IN    AL, DX
    AND   AL, 0FBh             ; clear bit 3
    OUT   DX, AL
    POP   DX
    JMP   rSW9

```

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```
rSW8:      ; enable window
           PUSH  DX
           MOV   DX, WINDOW_ENABLE_REG
           IN    AL, DX
           OR    AL, 04h           ; set bit 3
           OUT   DX, AL
           POP   DX

rSW9:      ; Check for valid access speed
           CMP   DL, 04           ; must be 100ns speed
           JE    rSW10
           MOV   AH, BadSpeed
           STC
           RET

rSW10:     XOR   AH, AH           ; Return success (resets [CF])
           RET

rSetWindow ENDP
```

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```

COMMENT ~*****~
Procedure:  rGetPage

    Entry:      [AL] = Adapter
                [BH] = Window
                [BL] = Page

    Exit:        [DL] = Page state
                Bit 0  Common space or attribute space
                Bit 1  Disabled or enabled
                [DI] = Memory card offset (4 KByte units)
*****~
PUBLIC rGetPage
rGetPage  PROC  NEAR
          ASSUME  CS:_TEXT, DS:DGROUP, ES:NOTHING

          ;[DL] = Page state
          MOV    DX, WINDOW_ENABLE_REG
          IN     AL, DX
          MOV    DL, 00h           ; !wp + disabled + common
          TEST   AL, 04h
          JZ    rGP4
          MOV    DL, 02h           ; set enabled bit
rGP4:
          ;[DI] = Memory card offset (4 KByte units)
          MOV    DI, MemCardOffset

          XOR    AH, AH           ; Return success (resets [CF])
          RET

rGetPage  ENDP

COMMENT ~*****~
Procedure:  rSetPage

    Entry:      [AL] = Adapter
                [BH] = Window
                [BL] = Page
                [DL] = Page state
                Bit 0  Common space or attribute space
                Bit 1  Disabled or enabled
                Bit 2  WP - Write protection
                [DI] = Memory card offset (4 KByte units)

    Exit:        All registers preserved
*****~
PUBLIC rSetPage
rSetPage  PROC  NEAR
          ASSUME  CS:_TEXT, DS:DGROUP, ES:NOTHING

          ;[DL] = Page state
          ; Check for valid page state
          CMP    DL, 00           ; common, disabled, !WP
          JE    rSP4
          CMP    DL, 02           ; common, enabled, !WP
          JE    rSP5
          MOV    AH, BadAttribute
          STC

          ;[DL] = Page state
          ; Check for valid page state
          CMP    DL, 00           ; common, disabled, !WP
          JE    rSP4
          CMP    DL, 02           ; common, enabled, !WP
          JE    rSP5
          MOV    AH, BadAttribute
          STC

```

```

                RET

rSP4:          ; disable window
                MOV     DX, WINDOW_ENABLE_REG
                IN      AL, DX
                AND     AL, 0FBh           ; clear bit 3
                OUT     DX, AL
                JMP     rSP6

rSP5:          ; enable window
                MOV     DX, WINDOW_ENABLE_REG
                IN      AL, DX
                OR      AL, 04h           ; set bit 3
                OUT     DX, AL

rSP6:          ;[DI] = Memory card offset (4 KByte units)
                ; 0) Save new offset in global variable
                MOV     MemCardOffset, DI

                ; 1) Determine window size and save multiplier in cl
                MOV     DX, WINDOW_SIZE_REG
                IN      AL, DX           ; a) read window size
                NOT     AL               ; b) complement
                AND     AL, 03h         ; c) mask off all but bits 0 and 1
                INC     AL               ; d) add 1
                MOV     CL, AL          ; e) save in cl (shift factor)
                MOV     CH, AL          ; f) save in ch (shift factor)

                ; 2) Calculate Page value, store in AL
                MOV     AX, DI
                XOR     AH, AH           ; highest value is 00FE
                SHR     AL, CL           ; c) Divide offset by WinSize/4K

                ; 3) Swap the page (AL) bits from low to high
                ;     i.e. 0000 0001 ==> 1000 0000
                ;     BL == i, AL == page, AH == fpage
                XOR     BL, BL           ; for loop counter i=0, for (i=0;;)

rSP7:          MOV     CL, BL
                MOV     DL, 1
                SHL     DL, CL           ; (1 << i)
                TEST    AL, DL           ; if ((page & (1 << i)) != 0)
                JZ      rSP8
                MOV     DL, 7
                SUB     DL, CL           ; (7 - i)
                MOV     CL, DL
                MOV     DL, 1
                SHL     DL, CL           ; 1 << (7 - i)
                OR      AH, DL           ; fpage |= 1 << (7 - i);

rSP8:          INC     BL               ; for (;i++)
                CMP     BL, 8           ; for (;i<8;)
                JB      rSP7

                ; 4) Output the right page number
                MOV     CL, CH           ; restore shift factor from l.f
                SHR     AH, CL           ; fpage >> i

```

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```

        MOV     AL, AH

        MOV     DX, WINDOW_PAGE_REG
        OUT    DX, AL

        XOR     AH, AH           ; Return success (resets [CF])
        RET
rSetPage ENDP

COMMENT ~*****~
Procedure:  rGetSocket

Entry:     [AL] = Adapter
          [BL] = Socket

Exit:     [BH] = Status change interrupt mask
          Bit 0  Write Protect Change
          Bit 1  Card Lock Change
          Bit 2  Ejection Request
          Bit 3  Insertion request
          Bit 4  Battery dead change
          Bit 5  Battery warning change
          Bit 6  Ready change
          Bit 7  Card detect change
          [CH] = VCC level
          [CL] = VPP1 & VPP2 levels
          [DH] = Socket state
          Bit 0  Write Protect Change
          Bit 1  Card Lock Change
          Bit 2  Ejection Request
          Bit 3  Insertion request
          Bit 4  Battery dead change
          Bit 5  Battery warning change
          Bit 6  Ready change
          Bit 7  Card detect change
          [DL] = Control and indicators state
          Bit 0  Write protect status
          Bit 1  Card lock status
          Bit 2  Motorized card ejection
          Bit 3  Motorized card ejection
          Bit 4  Card lock
          Bit 5  Battery status
          Bit 6  Busy status
          Bit 7  XIP status
          [DI] = IRQ level steering/Interface type

*****~
        PUBLIC rGetSocket
rGetSocket PROC NEAR
        ASSUME     CS:_TEXT, DS:DGROUP, ES:NOTHING
; [BH] = Status change interrupt mask
        MOV     BH, 0           ; none
; [CH] = VCC level
        MOV     CH, 0           ; no pwrentry
; [CL] = VPP1 & VPP2 levels
        MOV     CL, 0           ; no pwrentry
; [DH] = Socket state

```

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```

        MOV     DH, 0                ; no states to report
; [DL] = Control and indicators state
        MOV     DL, 0                ; none
; [DI] = LowByte=IRQ level steering/HiByte=Interface type
        MOV     DI, 0100h           ; No IRQ's, memory only
        XOR     AH, AH              ; Return success (resets [CF])
        RET
rGetSocket ENDP

COMMENT ~*****~
Procedure: rSetSocket

Entry:    [AL] = Adapter
          [BL] = Socket
Exit:     [BH] = Status change interrupt mask
          Bit 0 Write Protect Change
          Bit 1 Card Lock Change
          Bit 2 Ejection Request
          Bit 3 Insertion request
          Bit 4 Battery dead change
          Bit 5 Battery warning change
          Bit 6 Ready change
          Bit 7 Card detect change
          [CH] = VCC level
          [CL] = VPP1 & VPP2 levels
          [DH] = Socket state
          Bit 0 Write Protect Change
          Bit 1 Card Lock Change
          Bit 2 Ejection Request
          Bit 3 Insertion request
          Bit 4 Battery dead change
          Bit 5 Battery warning change
          Bit 6 Ready change
          Bit 7 Card detect change
          [DL] = Control and indicators state
          Bit 0 Write protect status
          Bit 1 Card lock status
          Bit 2 Motorized card ejection
          Bit 3 Motorized card ejection
          Bit 4 Card lock
          Bit 5 Battery status
          Bit 6 Busy status
          Bit 7 XIP status
          [DI] = IRQ level steering/Interface type

Exit:     All register preserved
*****~

PUBLIC rSetSocket
rSetSocket PROC NEAR
ASSUME    CS:_TEXT, DS:DGROUP, ES:NOTHING

        XOR     AH, AH              ; Return success (resets [CF])
        RET

rSetSocket ENDP

```

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```

COMMENT ~*****~

Procedure:  rGetStatus

Entry:      [AL] = Adapter
            [BL] = Socket

Exit:       [BH] = Current card state:
            Bit 0  Write protect
            Bit 1  Card Locked
            Bit 2  Ejection Request
            Bit 3  Insertion Request
            Bit 4  Battery dead
            Bit 5  Battery low
            Bit 6  Card ready
            Bit 7  Card detected
            [DH] = Socket state
            [DL] = Control and indicator state
            [DI] = IRQ level steering/Interface type

*****~

PUBLIC rGetStatus
rGetStatus PROC NEAR
ASSUME     CS:_TEXT, DS:DGROUP, ES:NOTHING

; [BH] = Status change interrupt mask
MOV  BH, 80h          ; card inserted

; [DH] = Socket state
MOV  DH, 0           ; no states to report

; [DL] = Control and indicators state
MOV  DL, 0           ; none

; [DI] = LowByte=IRQ level steering/HiByte=Interface type
MOV  DI, 0100h       ; No IRQ's, memory only

XOR  AH, AH          ; Return success (resets [CF])
RET

rGetStatus ENDP

_TEXT     ENDS

_INIT    SEGMENT BYTE PUBLIC 'INIT' USE16
ASSUME   CS:DGROUP, DS:DGROUP, ES:NOTHING

```

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```

COMMENT /*****
Procedure:   RfaInit

Purpose:    Do first-time initialization for RFA

Entry:      nothing

Exit:       [AX] = status

Initialize chip set and configure window for RFA
*****/
PUBLIC RfaInit
RfaInit PROC NEAR
ASSUME      CS:DGROUP, DS:DGROUP, ES:NOTHING

PUSH  BX
PUSH  DX

;          0) Setup REDWOOD Registers

CLI                      ; Must not be interrupted
IODELAY
MOV   DX, REDWOOD_INDEX
MOV   AX, REDWOOD_PS_REG3 ; 112h
OUT   DX, AX
IODELAY
MOV   DX, REDWOOD_DATA
IN    AX, DX
OR    AX, 4
OUT   DX, AX

IODELAY
MOV   DX, REDWOOD_INDEX
MOV   AX, REDWOOD_PS_REG1 ; 110h
OUT   DX, AX
IODELAY
MOV   DX, REDWOOD_DATA
IN    AX, DX
OR    AX, 0100h
OUT   DX, AX
STI

;          1) Disable FlashROMdisk (RFA)
MOV   DX, WINDOW_ENABLE_REG
IN    AL, DX
AND   AL, 0FBh          ; clear bit 3
OUT   DX, AL

;          2) Configure system window base address
MOV   BX, rMinWinBase
call  SetWindowBaseAddr

;          3) Configure system window size
MOV   DX, WINDOW_SIZE_REG
IN    AL, DX
AND   AL, 0FCh          ; Clear bits 0,1
OR    AL, WINDOW_SIZE_16K ; assign bits 0,1

```

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```
        OUT    DX, AL

;        4) Configure page address (card address 0)
MOV     MemCardOffset, 0
MOV     AL, 0
MOV     DX, WINDOW_PAGE_REG
OUT     DX, AL

;        5) Enable FlashROMdisk (RFA)
MOV     DX, WINDOW_ENABLE_REG
IN      AL, DX
OR      AL, 04h           ; set bit 3
OUT     DX, AL

POP     DX
POP     BX

XOR     AX, AX           ; Return success (resets [CF])
RET

RfaInit ENDP

_INIT   ENDS

END
```

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